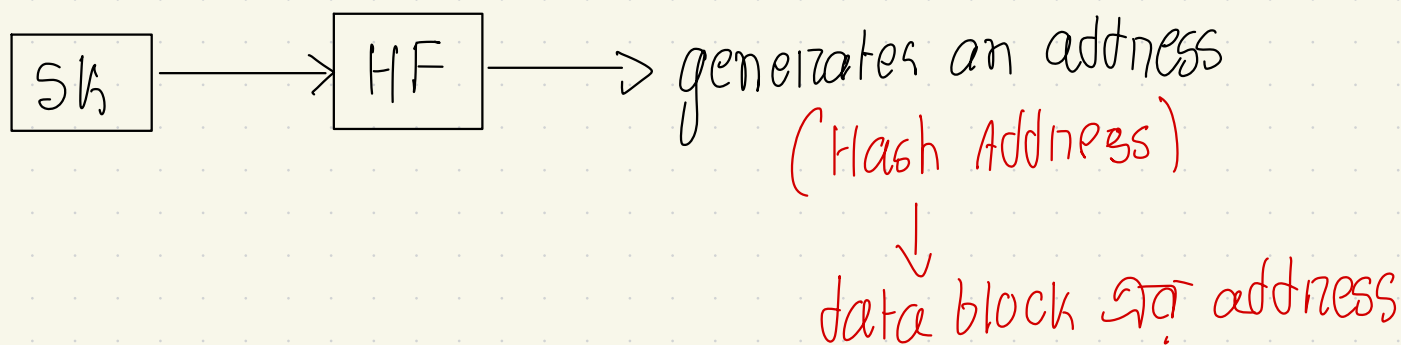


24/08/25

# Static Hashing



Hashing  $\rightarrow$  no separate index file

SK	HA

$\hookrightarrow$  optional  
(no need of indexing)

⊛ Hashing by mul } hash function  
⊛ " " div }

$$\text{st id} = 10 = k$$

$$k \bmod 3$$

$$10 \div 3 = \textcircled{1}$$

[ st. id = 10 stored in data block 1 ]

Sk = CAT

static  $\rightarrow$  no. of bucket and size will be fixed

Bucket = 10

$$C = 3$$

$$A = 1$$

$$T = 20$$

$$Sk = 24$$

$$24 \bmod 10$$

$$= 4$$

(stored in no. 4 bucket)

Sk at 2nd bucket and 2nd

secondary index file is dense



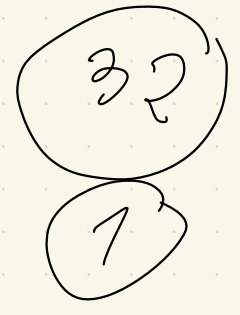
SK

H/A

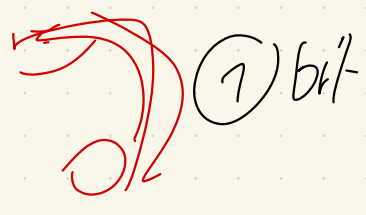
1	11010
2	00000
3	1110
4	00000
5	01001
6	10101
7	10111



$i = 32$   
 $i < 32$   
 $i > 1$

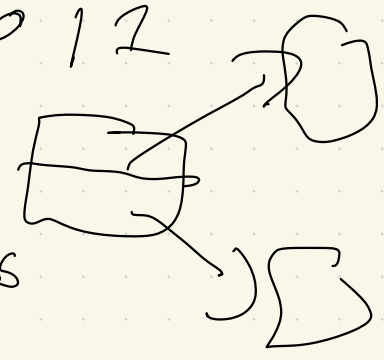


10110



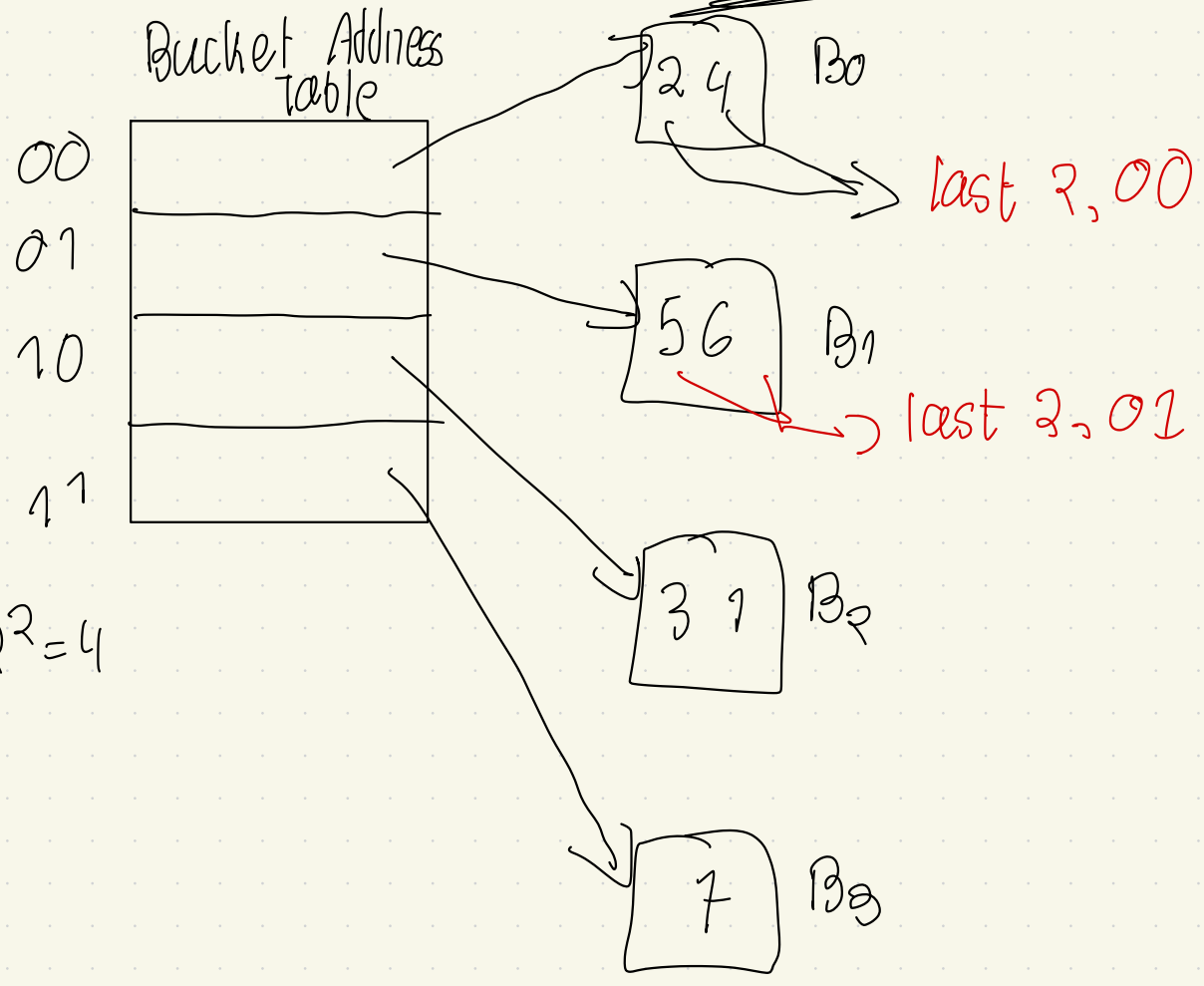
$2^i$

0 1 2



$2^i \rightarrow$  rows in bucket table  
dynamic hashing can generate 32 bit address

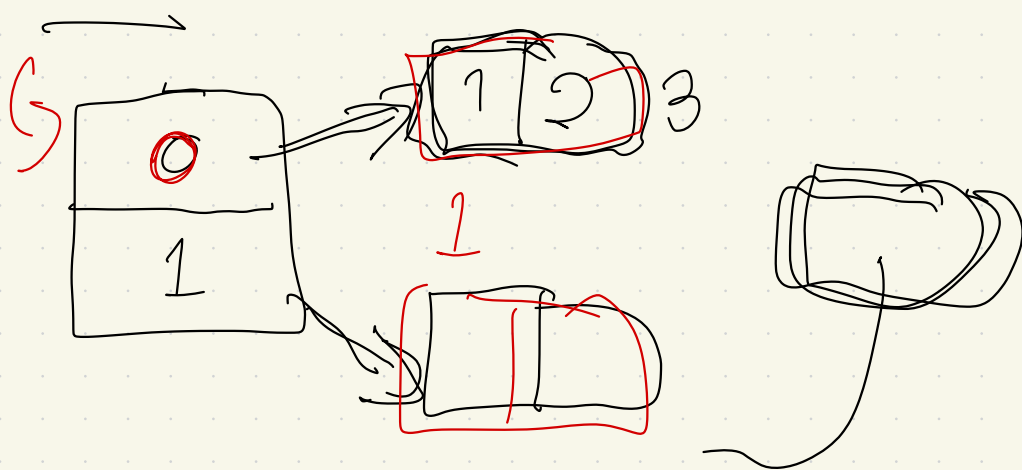
max 2 values



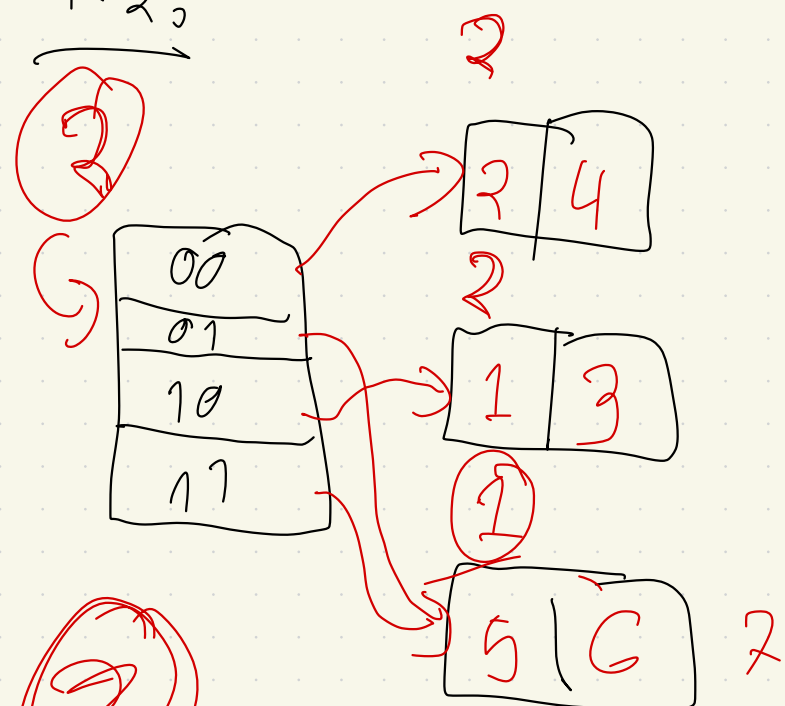
SK = 2, H/A = 10001

but no space to add also if I add another bucket, no table now to link the bucket.

$i = 1$

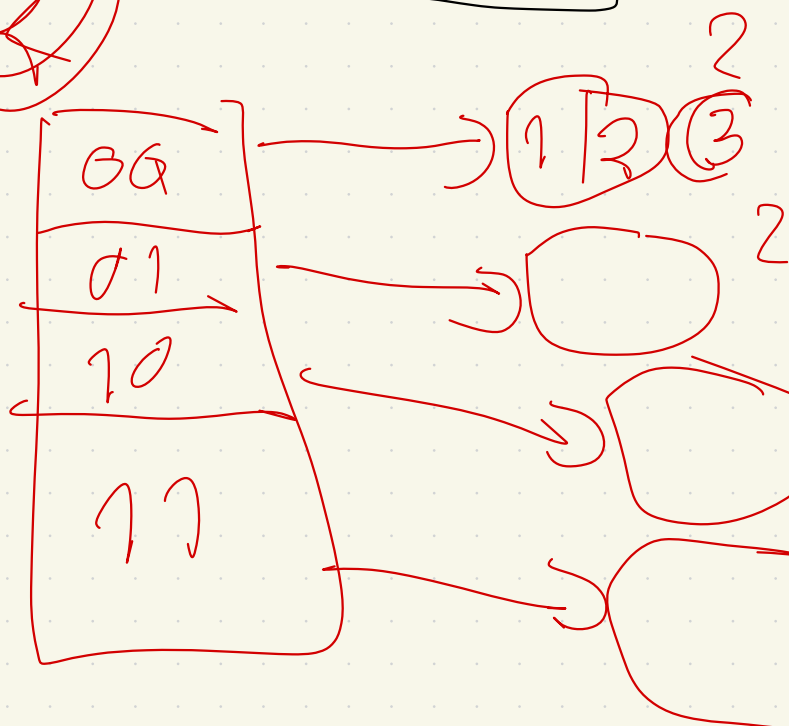


$i = 2$



00  
10

01  
11

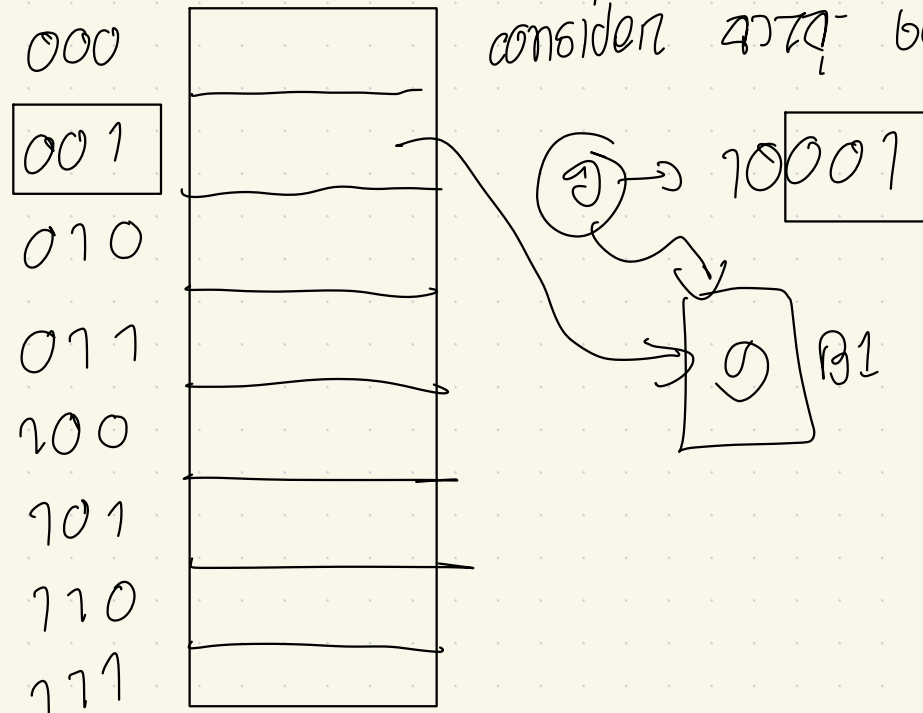


2

Hence,  
now we will increase the value of  $i$  which we  
took 2 initially.  $i=3$   $2^3=8$

now show value of last 3 bits

consider bucket 1 and bucket



12 \* 64 slides  $\rightarrow$  next am

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$T_i$	$T_j$
$I_i$	$I_j(Q)$
$I_i$	$I_j(Q)$

conflict:

- ① write operation
- ② same data

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one schedule is serial and another is concurrent  
They are equivalent if they have same no. of  
instructions and if they are consistent

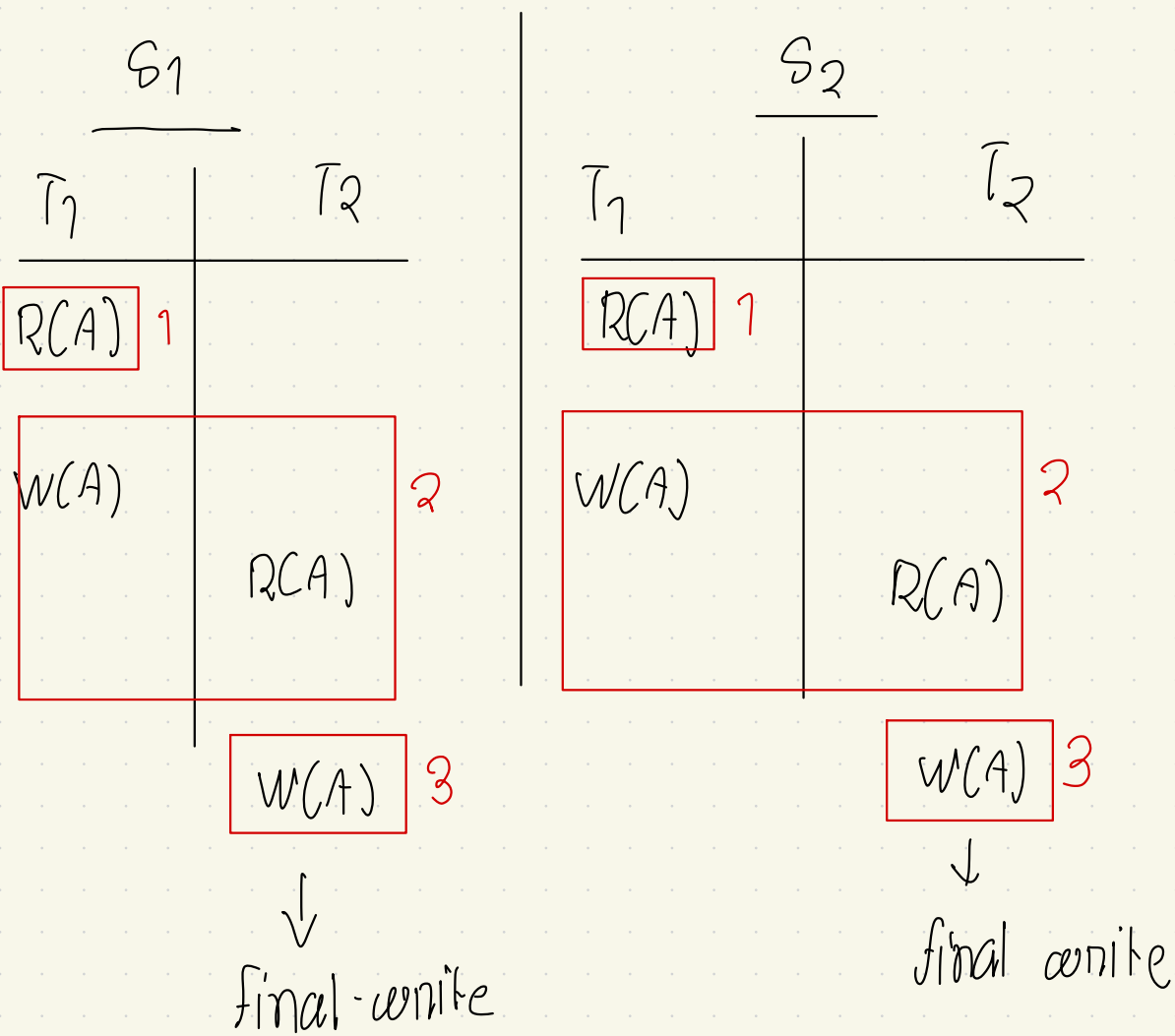
\* Serial is always consistent

Consistency check

एक schedule एक serial बनाने में सक्षम है like  
equivalent to serial schedule बनाने में सक्षम है then  
can be made consistent.

\* Same data read, write after write access conflict multiple read will be no conflict.  
 conflict equivalent  $\rightarrow$  swap non-conflicting commands.

## # View Serializability:



equivalent  $\rightarrow$  all three conditions must fulfill  
 Agar condition ka access kahi condition match  
 karne view equivalent.

if a transaction starts with write operation, that is called blind write. Those schedule are if view equivalent but not conflict equivalent.

⑧ previous stable state is previously commit state.

⑧ Cascaded roll back.

↳ problem if too many transactions. Waste of time.

02/09/25

stable state  $\rightarrow$  commit

ଆଉ ଓହ୍ଲା dependent କରି ଆଉ commit କରା

then read



T8	T9
read(A)	
write(A)	
	read(A)
write(A)	
commit	
	commit

Fig: Schedule-11

recoverable but T9 is  
 read(A) from data read by  
 T8. So T9 may  
 roll back after it

Q: What is cascading rollback?  
 write operation is not dependent.  
 cascadeless schedule is not serial schedule