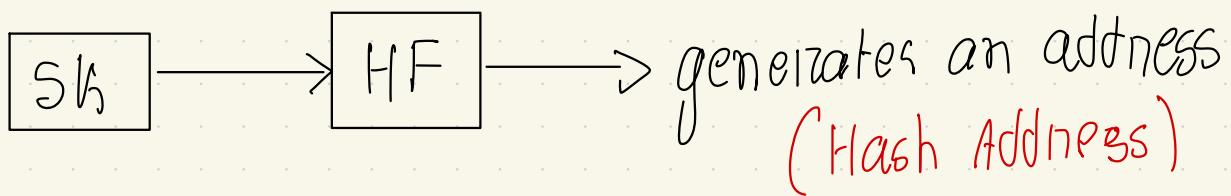


24/08/25

Static Hashing



↓
data block at address

Hashing ↗ no separate index file

SK	HA

↳ optional
(no need of indexing)

* Hashing by mul } hash function
* || . . . || div }

$$\text{st. id} = 10 = k$$

$$k \bmod 3$$

$$10 \div 3 = 1$$

[st. id = 10 stored in data block 1]

$Sk = CAT$ static \rightarrow no. of bucket and size will be fixed

Bucket = 10

$$C = 3$$

$$A = 1$$

$$T = 20$$

$$Sk = 24$$

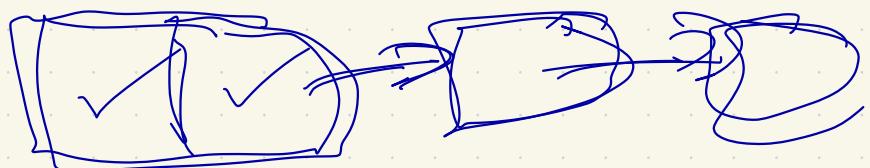
$24 \bmod 10$

$$= 4$$

(stored in no. 4 bucket)

Sk at 2nd bucket and 2nd

secondary index file is dense



SK

HA

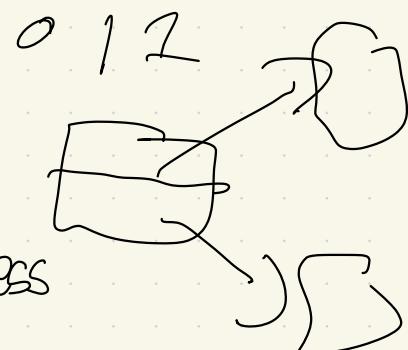
1	11010
2	00000
3	1110
4	00000
5	01001
6	10101
7	10111



$i = 32$
 ~~$i = 32$~~
 $i > 1$

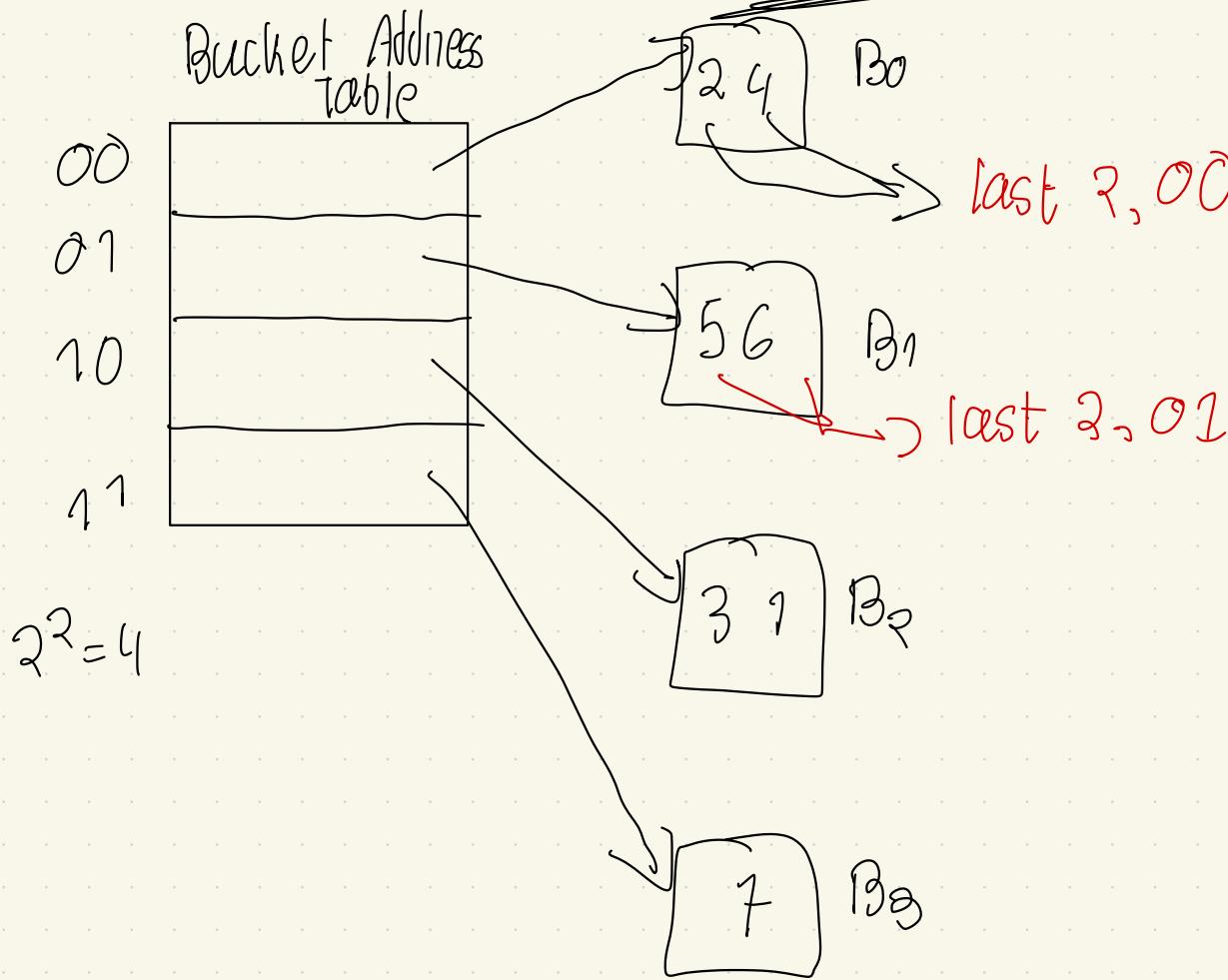
32
 1

10110
 ① bit
 2ⁱ



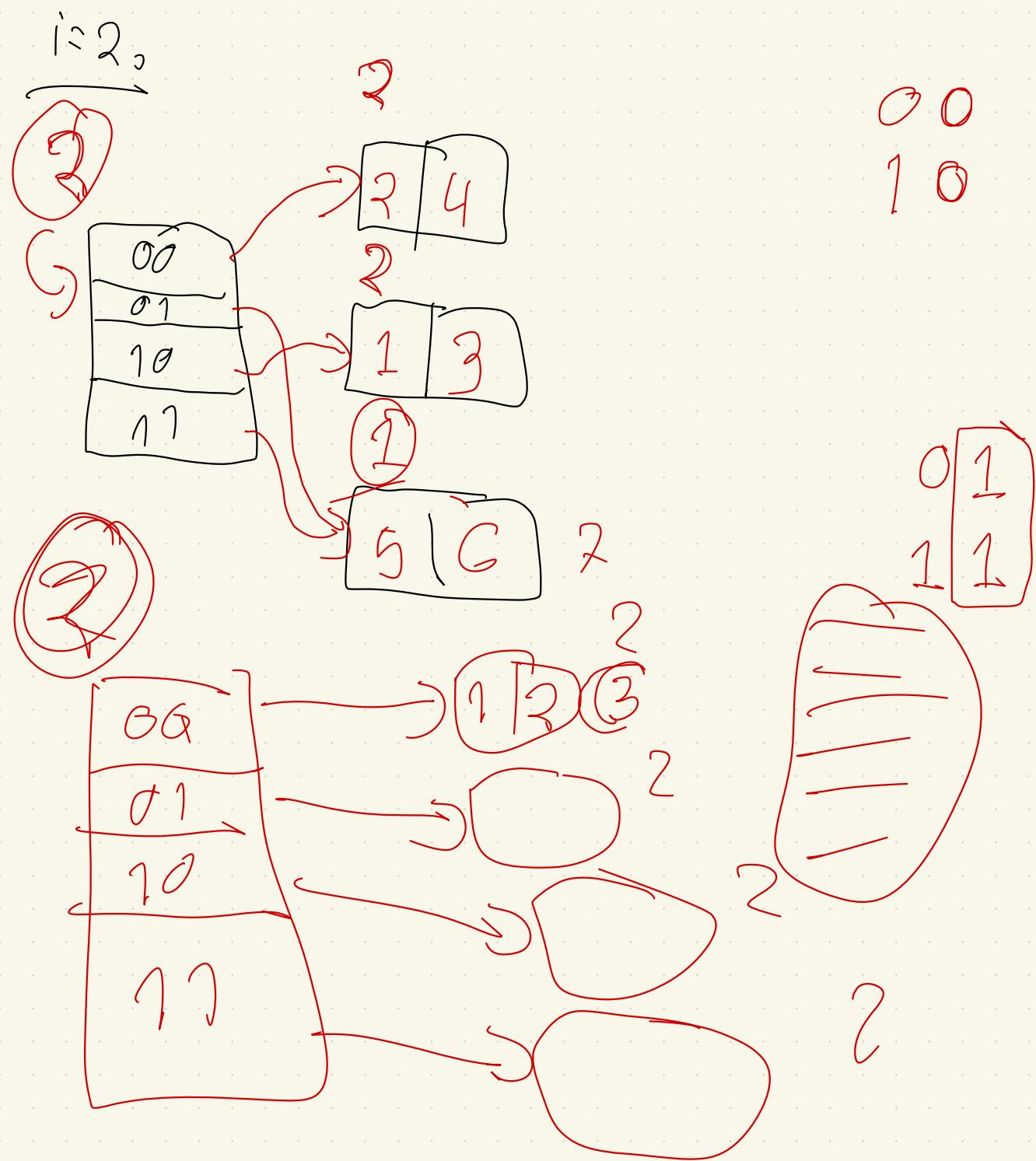
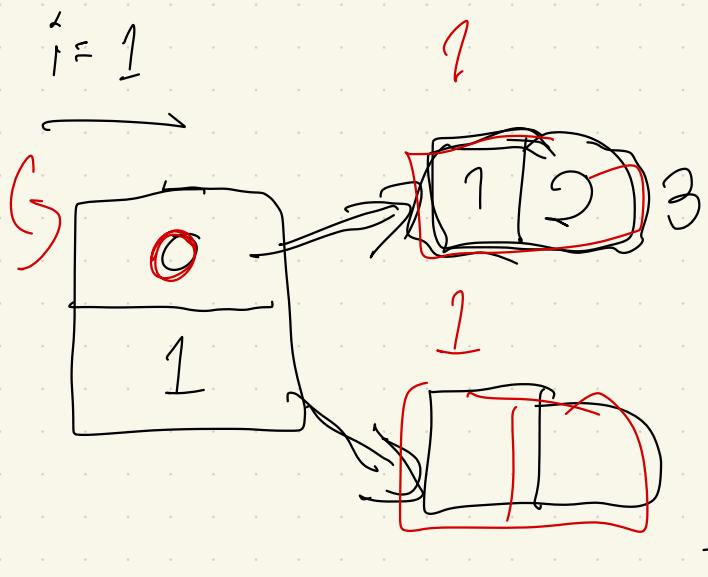
$2^i \rightarrow$ rows in bucket table
 dynamic hashing can generate 32 bit address

max 2 values



SK = 0, HA = 10001

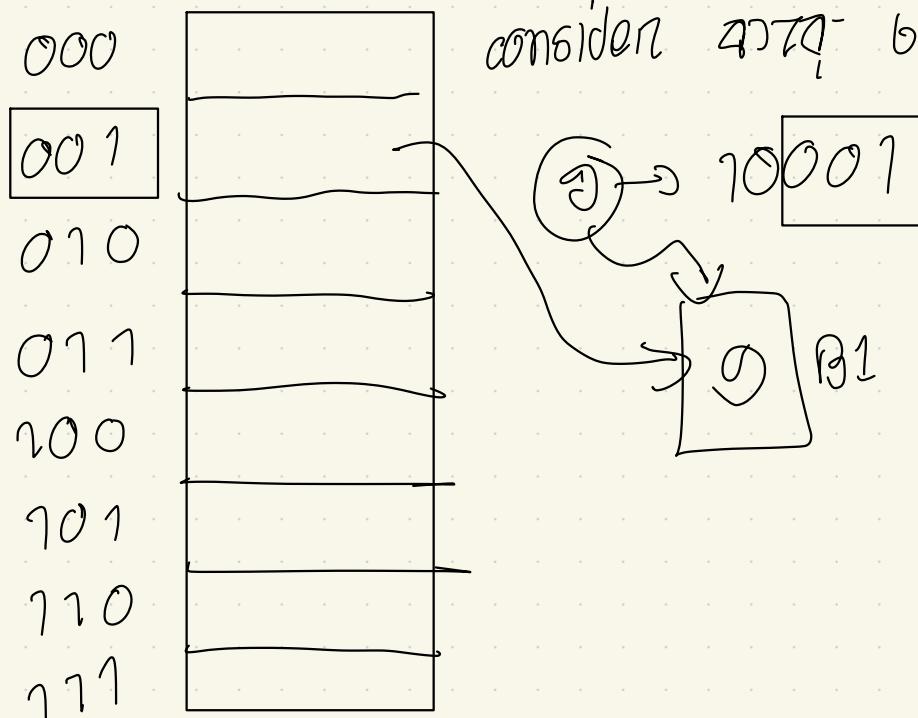
but no space to add also if I add another bucket, no table now to link the bucket.



Hence,
now we will increase the value of i which we
took 2 initially. $i=3 \quad 2^3=8$

now show value 610101- last 3 bits

consider 001- bucket 0 cannot



12 * 64 slides \rightarrow 1121 test am

26/08/25

T_i	T_j
I_i	$I_j(Q)$
I_i	$I_j(Q)$

conflict:

- ① write operation
- ② same data

31/08/25

one schedule is serial and another is concurrent

They are equivalent if they have same no. of instructions and if they are consistent

* Serial is always consistent

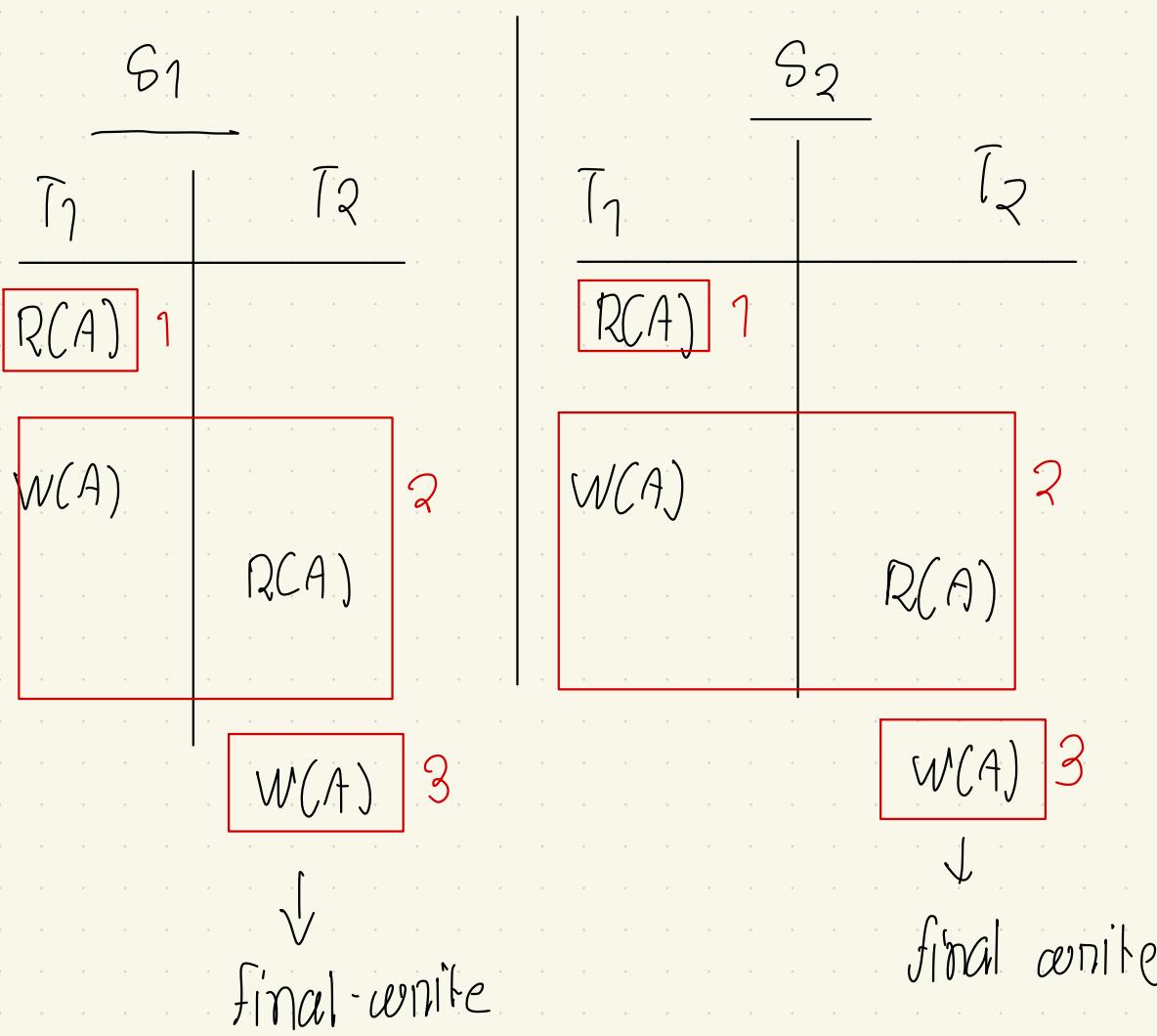
Consistency check

Two schedule that serial pattern appear like equivalent to serial schedule appear appear that will be made consistent.

④ Same data \Rightarrow no conflict write as it is conflict
multiple read will be no conflict.

conflict equivalent \rightarrow swap non-conflicting commands.

View Serializability:



equivalent \rightarrow all three conditions must fulfill
कोरा condition एवं दूसरी backi condition match
काढ़ाता view equivalent.

if a transaction starts with write operation, that is called blind write. Those schedule are if view equivalent but not conflict equivalent.

- ④ previous stable state is previously commit state.
- ⑤ cascaded roll back:
 - ↳ problem if too many transactions. Waste of time.

02/09/25

stable state → commit

→ - don't dependent on other commit and

then read

T8	T9
read(A)	
write(A)	
	read(A)
write(A)	
commit	
	commit

Fig: Schedule-11

recoverable but T9 is
 read(A) 훔은 data read from
 other transaction. So T9 fails
 roll back and T9 fails

Q: What is cascading rollback?
 write operation is not dependent.
 cascaded schedule is not serial schedule