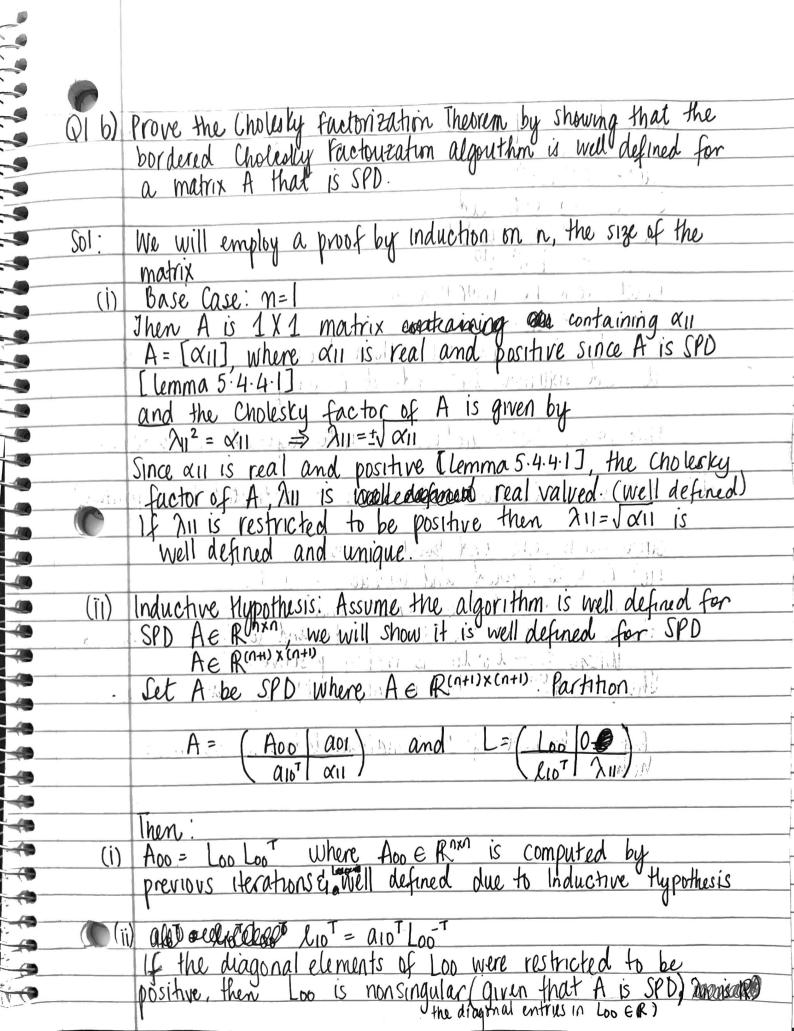
Exama Q1: Propose a bordered algorithm for computing the Cholesky factorization of a SPD matrix A. Sol. a) Consider A = LL Where L is a lower triangular matrix and A and L are partitioned as follows:
Exama OI: Propose a bordered algorithm for computing the Cholesky
factorization of a SPD matrix A. factorization of a SPD matrix A.
factorization of where Lis a lower triangular matrix and
factorization of a SPD matrix A. Sol: a) Consider A = LLT where L is a lower triangular matrix and A and L are partitioned as follows:
$A = \begin{pmatrix} Aoo & Qol \\ QoT & QUI \end{pmatrix} \text{ and } L = \begin{pmatrix} Loo & Q \\ LloT & \lambda II \end{pmatrix}$
$a_{io}^{T} \alpha_{ii} \rangle$
By substituting these partitioned matrices into A = LLT, we find
that:
(A00 a01) = (L00 D) (L00 D)
aid ail lot sil lot sil
= (Lob D) (Lob Lib) <transpose l="" of=""></transpose>
(lot AII) (b AII) (matrix matrix
So, Hoo wor (Plat I hat I hat All All) multiplication >
· Since A is SPD, we don't need to compute any (it is just a
transpose of air, Hence;
transpose of a10T), Hence; (A00 * = (L00 L00 * L00 * L10T L10T L10T L10T L10T L10T L10T L10T
· Hence
Hence $(i) Aoo = Loo Loo^{T} = Cholesky factorization of submatrix Aoo$ $(ii) \alpha_{10}^{T} = L_{10}^{T} Loo^{T}$ $(iii) \alpha_{11} = L_{10}^{T} L_{10} + \lambda_{11}^{Z}$
$(ii) \alpha_{10}^{T} = \mathcal{L}_{10}^{T} \mathcal{L}_{00}^{T}$ $(iii) \alpha_{11} = \mathcal{L}_{10}^{T} \mathcal{L}_{10} + \alpha_{11}^{T}$
(i) Assuming that Loo has been computed in previous iterations
1 1 - 1 on los' = (Innestitute Attorne De 100 can be obtained
by previous iterations of the algorithm
(ii) and = liot Loot
Assuming we know what allo' appelled (given) and Loo' computed
previously) is known, we can get his its follows.
Assuming we know what and anot analled (given) and Loo (computed previously) is known, we can get his as follows: \[\begin{align*} \text{Previously} & Previou

 $\propto 11 = l_{10}^{T} l_{10} + \lambda 11^{2}$ (where $(l_{10})^{T} = l_{10}^{T}$) (iii) Since we know an (given) and lipt (computed in step ii),

We can get liotlio; to get his; we can do the following: $\lambda_{11} = \sqrt{\alpha_{11} - l_{10}tl_{10}}$ So, the algorithm looks like: Loverwrites Acuse lower triangular A:= CHOLFACT_BORDERED (A) part of A w/ its Cholesky fac-Partition A -> (ATL * Where ATL is OXD matrix while n(ATL) < n(A) do late Repartition ATL X A00 1 ABL | ABR) ant an * Azo al Azal Where an is 1x1 matrix 1 116 land alot: = liot: = alot ADOT < computing who only lower triangular part αιι := λιι := \ αιι - 10000 αιο αιο αιο Continue with: A00 \$ A. alo all * 1 Apo 1001 A22/ enduhile missione and a least of har treated one for the



then its transpose Loo is also non-singular & its inverse LooT exist. Iransposing both sides
(lio^T) | T = (aio^T Loo^T) | T |
lio = Loo^T aio which can be rewritten as nLoo lion= a10 Since Loo is non singular, then Loo lio = aio has a well defined solytion lio and it is unique.

So, is (Loo lio) = (aio) > lio Loo = aio => lio = aio Loo Therefore: lio = aio Loo is well defined and unique. and other your last a mi All = Vall- liotlio (since All is restricted to be positive) to be listly of of Since all is real and positive and he had be liot is well defined and unique, so is lio then authorlio exists and is unique Hence, Au = Vau - liot lio exists and uniquely determined:

When an - liot lio is real & positive:

Thus, Ae Rinto xinto has a unique Chioestay factorization. By the Principle of Mathematical Induction, the result holds.

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