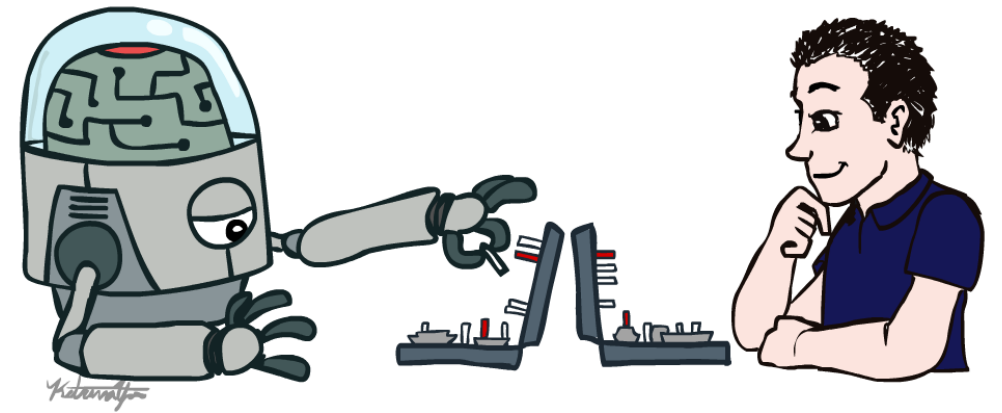


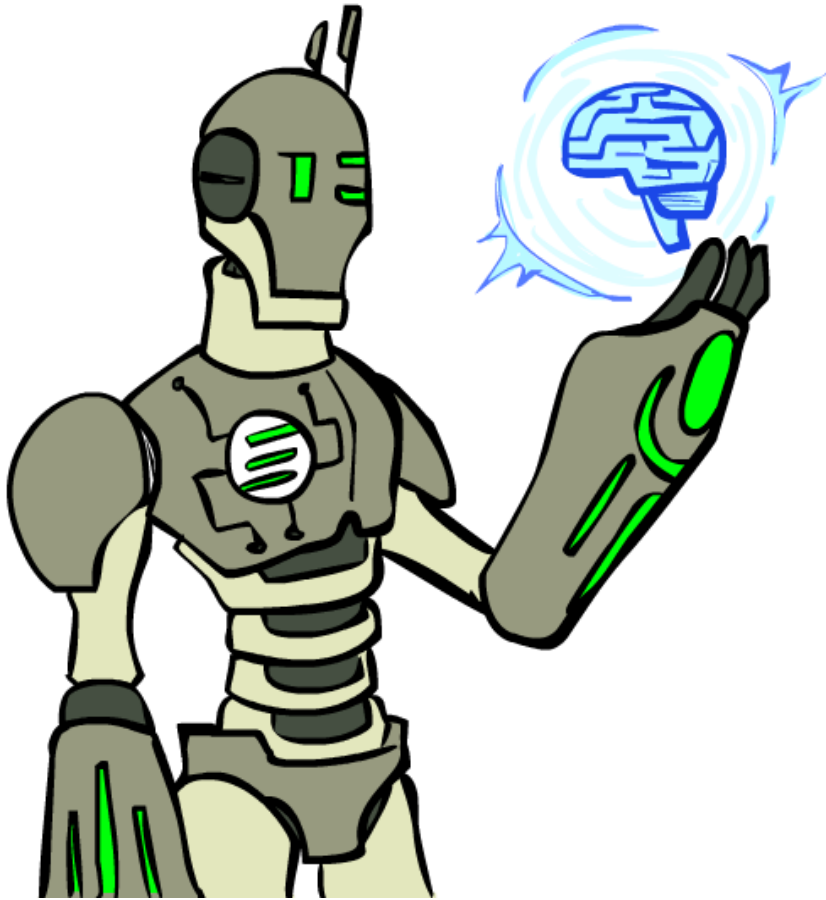
# Lecture 02

Ashis Kumar Chanda  
[chanda@rowan.edu](mailto:chanda@rowan.edu)



# Today

---



- Problem-solving agents
- Problem types
- Problem formulation
- Example problems
- Basic search algorithms

# Problem-solving agents

---

Function Sim\_Prob\_solv\_agent (*percept*) return an *action*

static: *seq*, an action sequence, initially empty

*state*, some description of the current world state

*goal*, a goal, initially null (based on current situation and PM)

*problem*, a problem formulation (what action and state to consider, given a goal)

*state*  $\leftarrow$  Update\_State (*state*, *percept*)

If *seq* is empty then do

*goal*  $\leftarrow$  Formulate\_Goal (*state*)

*problem*  $\leftarrow$  Formulate\_problem (*state*, *goal*)

*seq*  $\leftarrow$  Search (*problem*)

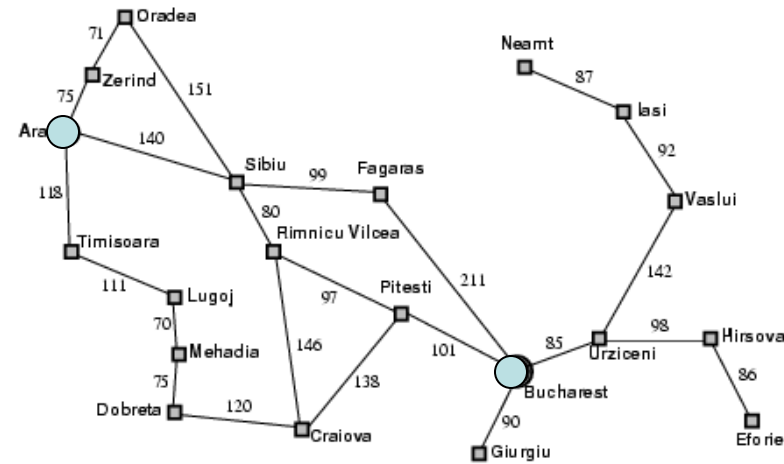
*action*  $\leftarrow$  First (*seq*)

*seq*  $\leftarrow$  Rest (*seq*)

Return *action*

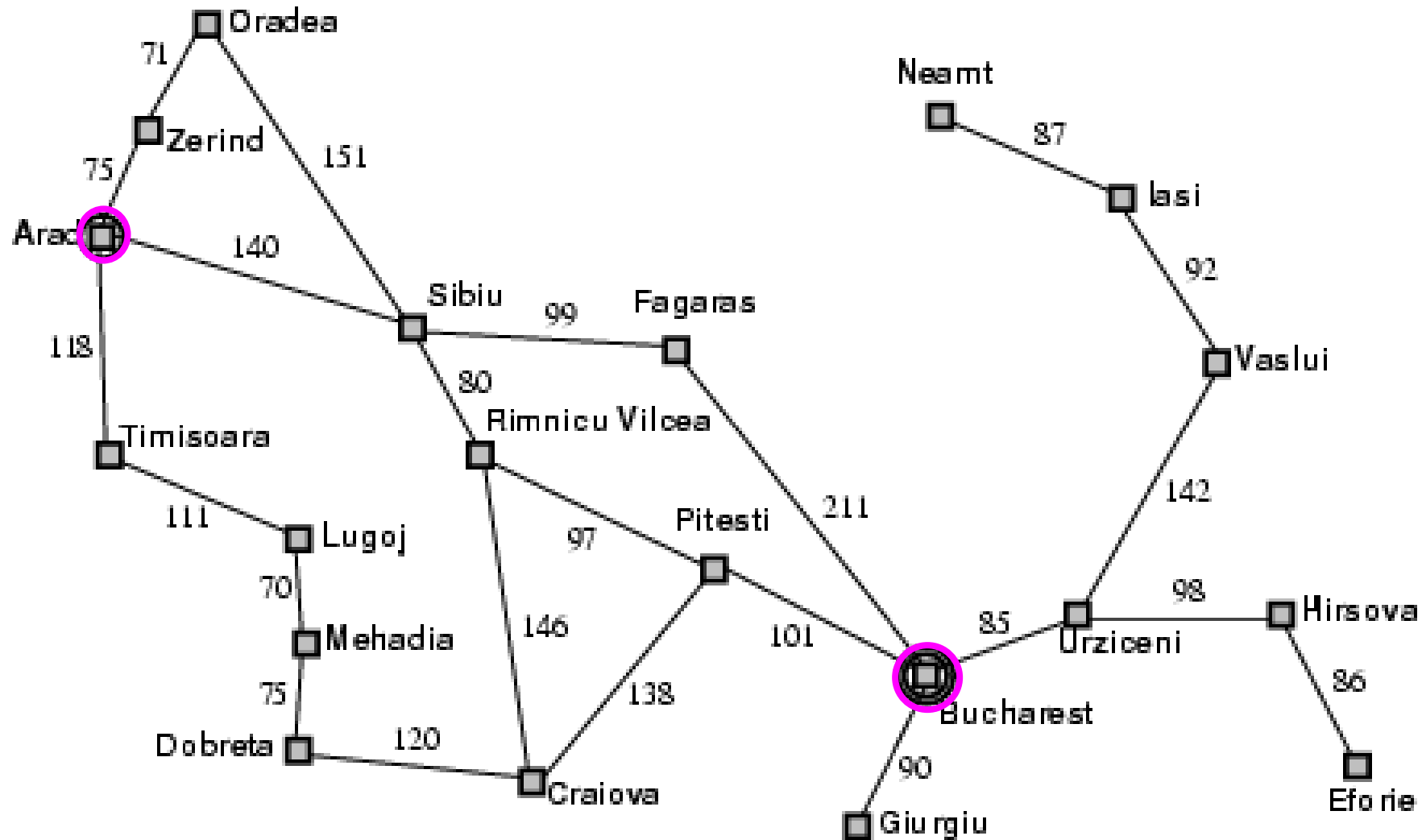
# Example: Romania

- Romania is a southeastern European country.
- Suppose, on holiday in Romania; you are currently in Arad.
- Flight leaves tomorrow from Bucharest
- Formulate goal:
  - be in Bucharest
- Formulate problem:
  - states: various cities
  - actions: drive between cities
- Find solution:
  - sequence of cities, e.g., Arad, Sibiu, Fagaras, Bucharest



function SEARCH (*problem*) return *solution* (as a sequence)

# Example: Romania



# Search problem formulation

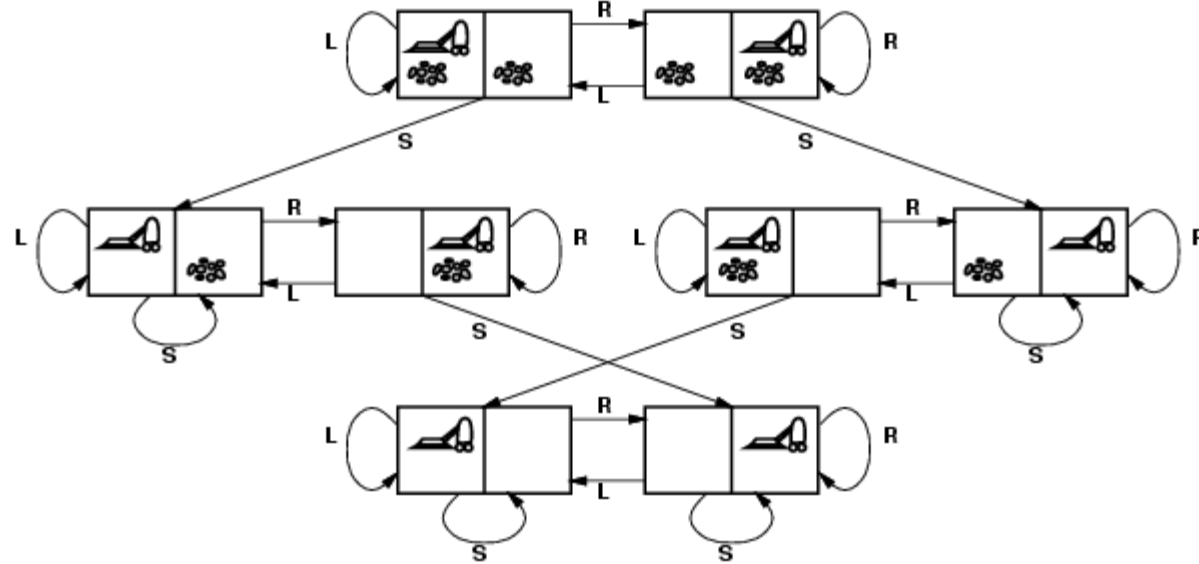
---

A **problem** is defined by four items:

1. initial state e.g., "at Arad"
2. **actions** or successor function  $S(x)$  = set of action–state pairs
  - e.g.,  $S(\text{Arad}) = \{ \langle \text{Arad} \rightarrow \text{Zerind}, \text{Zerind} \rangle, \dots \}$
3. **goal test**, can be
  - explicit, e.g.,  $x = \text{"at Bucharest"}$
  - implicit, e.g.,  $\text{Checkmate}(x)$
4. **path cost**
  - e.g., sum of distances, number of actions executed, etc.
  - $c(x, a, y)$  is the step cost, assumed to be  $\geq 0$

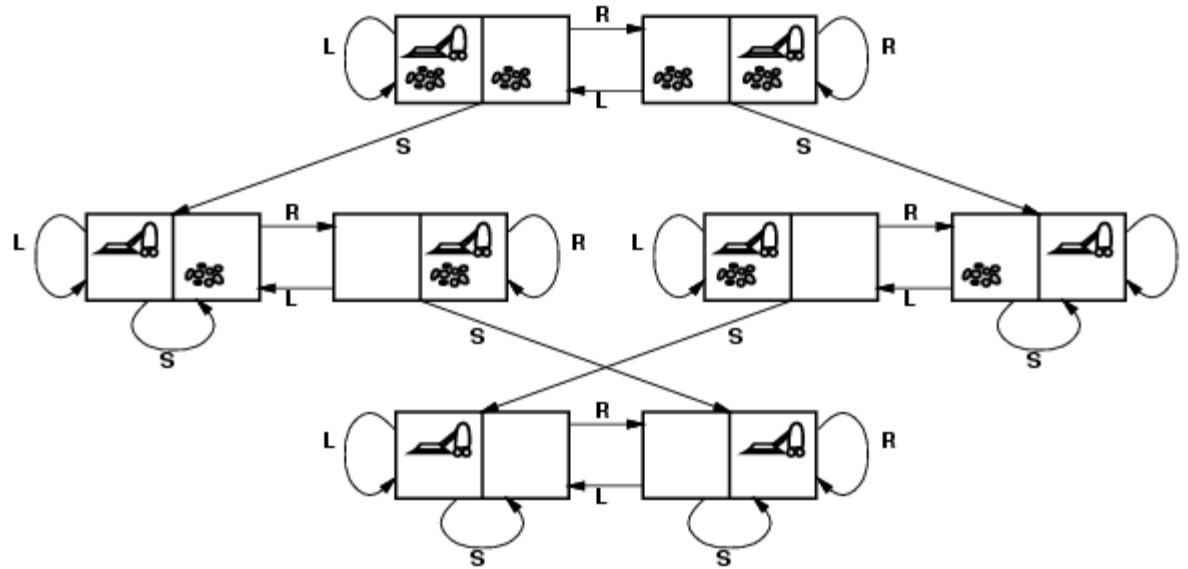
A **solution** is a sequence of actions leading from the initial state to a goal state

# Vacuum world state space graph



- states?
- actions?
- goal test?
- path cost?

# Vacuum world state space graph



- states? dirt and robot location
- actions? *Left, Right, Clean*
- goal test? no dirt at all locations
- path cost? 1 per action



# Example: The 8-puzzle

---

7	2	4
5		6
8	3	1

Start State

	1	2
3	4	5
6	7	8

Goal State

- states?
- actions?
- goal test?
- path cost?

# Example: The 8-puzzle

---

7	2	4
5		6
8	3	1

Start State

	1	2
3	4	5
6	7	8

Goal State

- states? locations of tiles
- actions? move blank left, right, up, down
- goal test? = goal state (given)
- path cost? 1 per move

# Tree search algorithms

---

- Basic idea:
  - offline, simulated exploration of state space by generating successors of already-explored states.

Function Tree\_search (*problem*, *strategy*) returns a *solution* or *failure*

    Initialize the search tree using the initial state of *problem*

    Loop do

        If there are no candidate for expansion then return *failure*

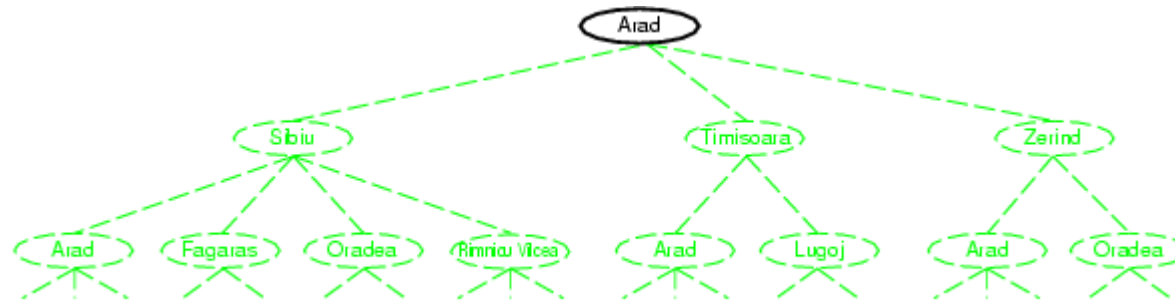
        Choose a leaf node for expansion according to *strategy*

        if the node contain a goal state then return the corresponding solution

        else expand the node and add the resulting nodes to the search tree

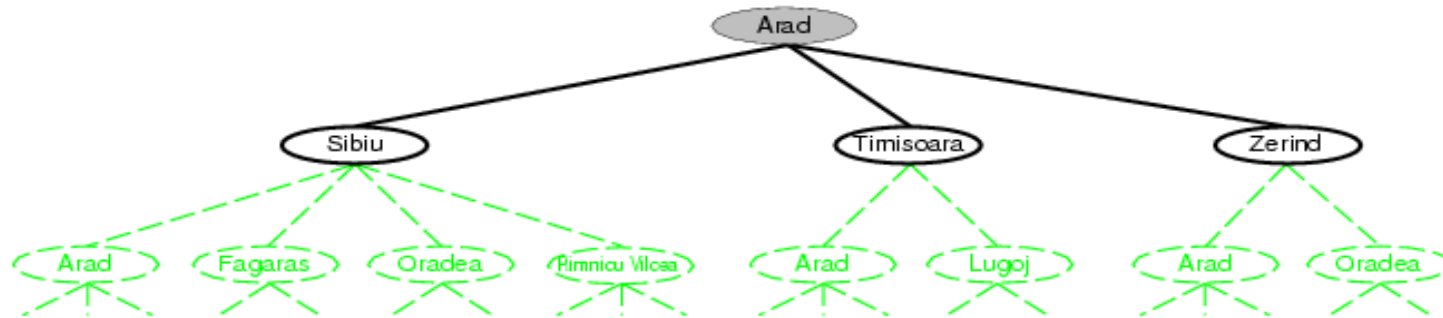
# Tree search example

---



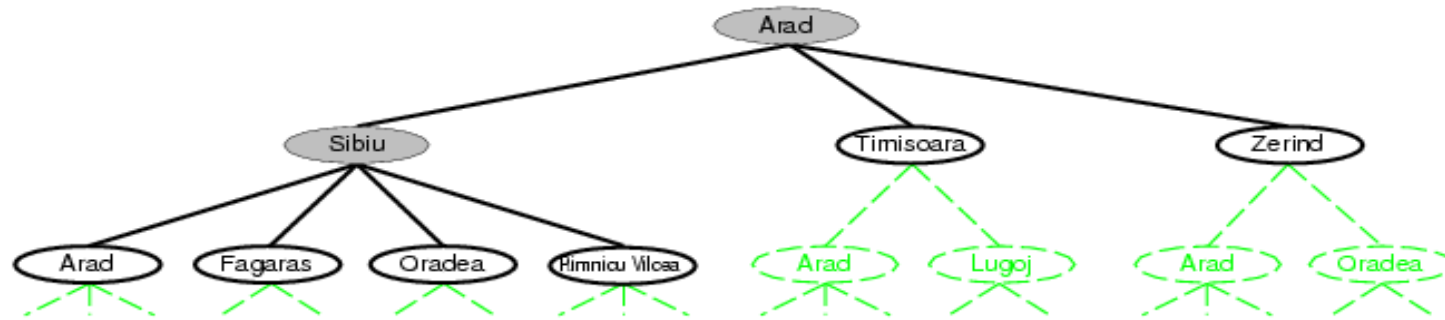
# Tree search example

---



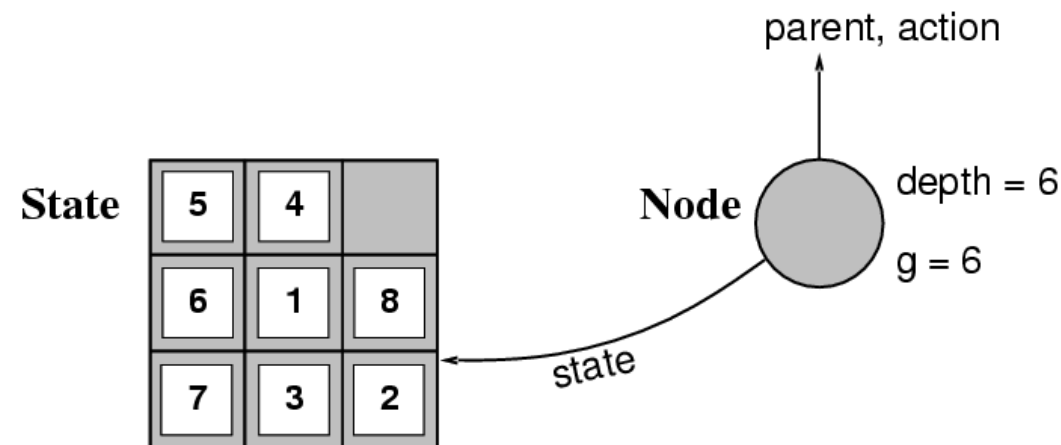
# Tree search example

---



# Implementation: states vs. nodes

- A **state** is a (representation of) a physical configuration
- A **node** is a data structure constituting part of a search tree includes **state**, **parent node**, **action**, **path cost  $g(x)$** , **depth**



# Search strategies

---

- A **search strategy** is defined by picking the **order of node expansion**
- Strategies are evaluated along the following dimensions:
  - **completeness**: does it always find a solution if one exists?
  - **time complexity**: number of nodes generated
  - **space complexity**: maximum number of nodes in memory
  - **optimality**: does it always find a least-cost solution?
- Time and space complexity are measured in terms of
  - ***b***: maximum branching factor of the search tree
  - ***d***: depth of the **least-cost** solution
  - ***m***: maximum depth of the state space (may be  $\infty$ )



# Uninformed search strategies

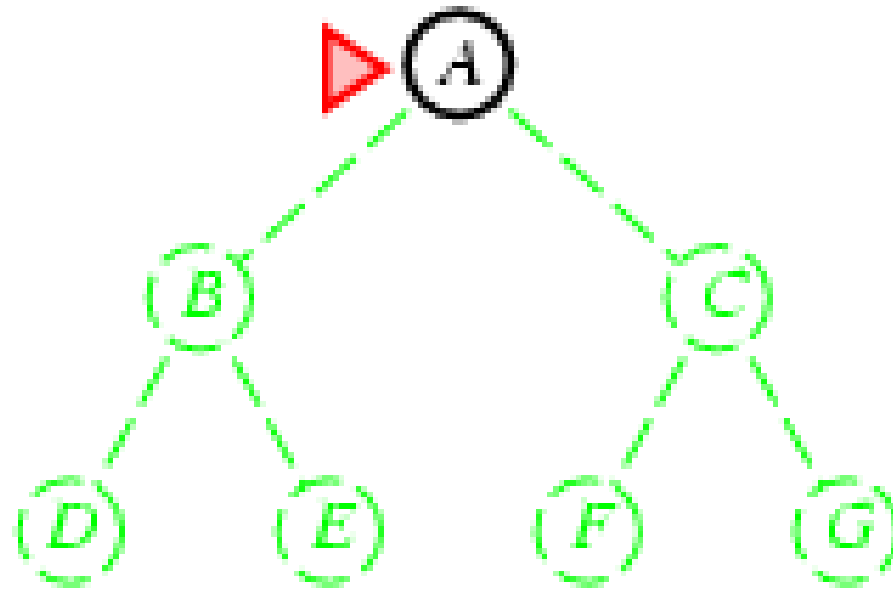
---

- **Uninformed** (**blind**) search strategies use only the information available in the problem definition
- Breadth-first search
- Uniform-cost search
- Depth-first search
- Depth-limited search
- Iterative deepening search

# Breadth-first search

---

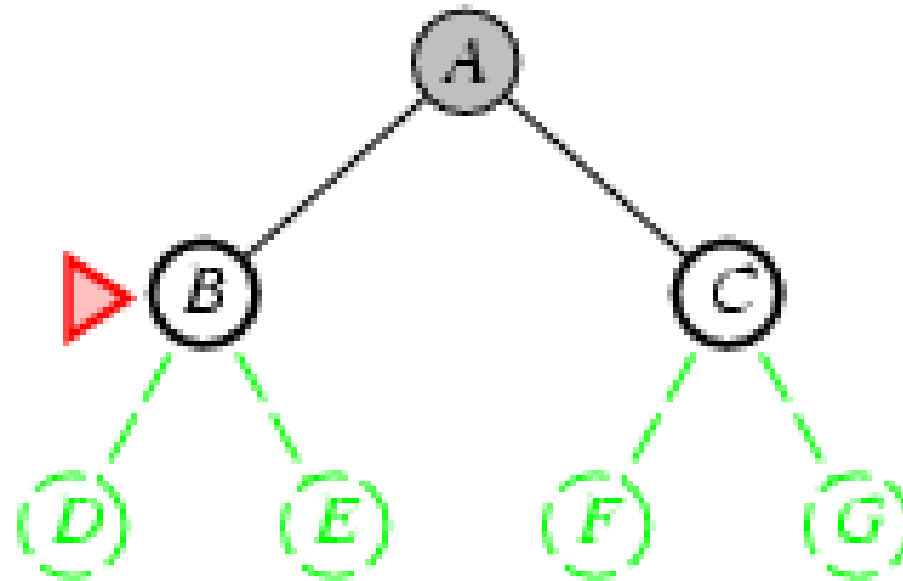
- Expand shallowest unexpanded node
- **Implementation:**
  - *fringe* is a FIFO queue, i.e., new successors go at end



# Breadth-first search

---

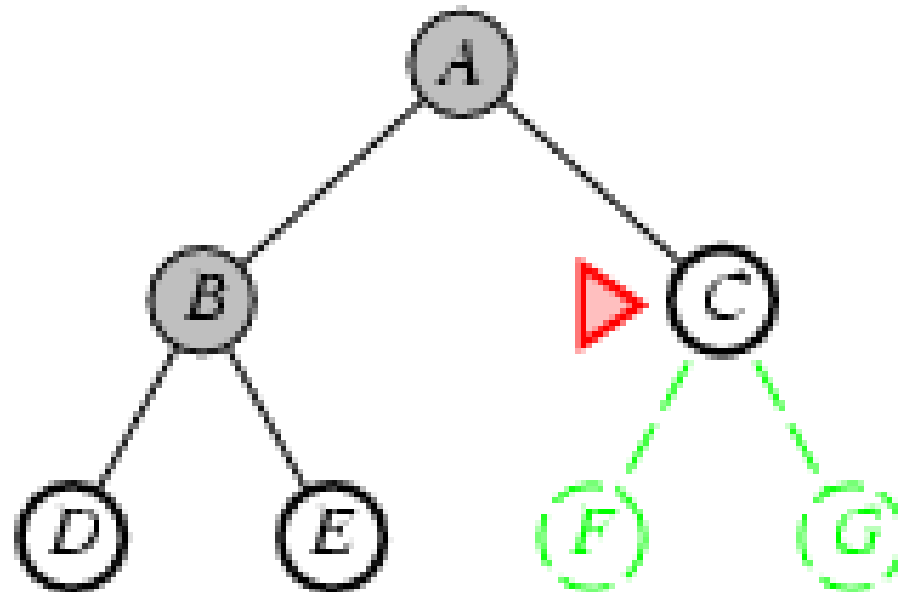
- Expand shallowest unexpanded node
- **Implementation:**
  - *fringe* is a FIFO queue, i.e., new successors go at end



# Breadth-first search

---

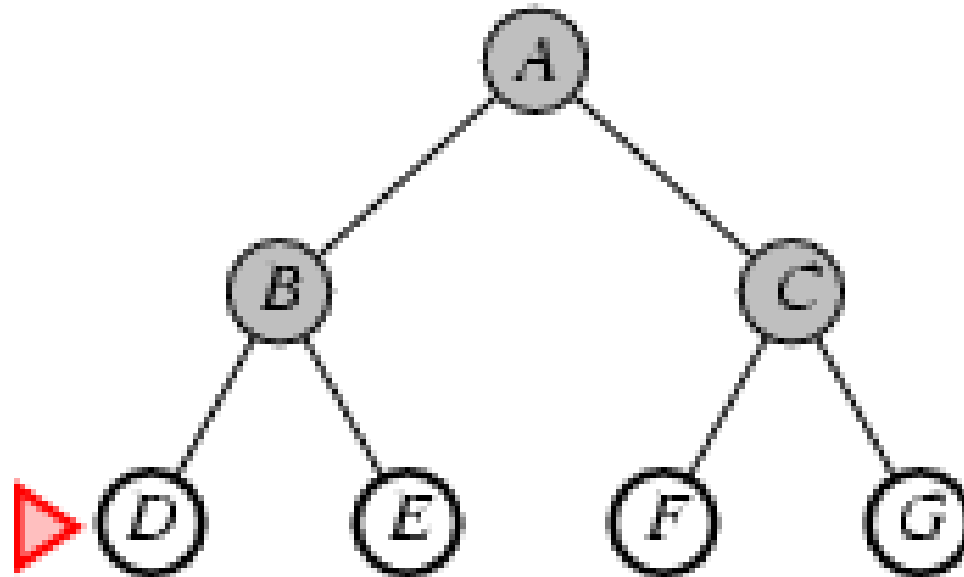
- Expand shallowest unexpanded node
- **Implementation:**
  - *fringe* is a FIFO queue, i.e., new successors go at end



# Breadth-first search

---

- Expand shallowest unexpanded node
- **Implementation:**
  - *fringe* is a FIFO queue, i.e., new successors go at end



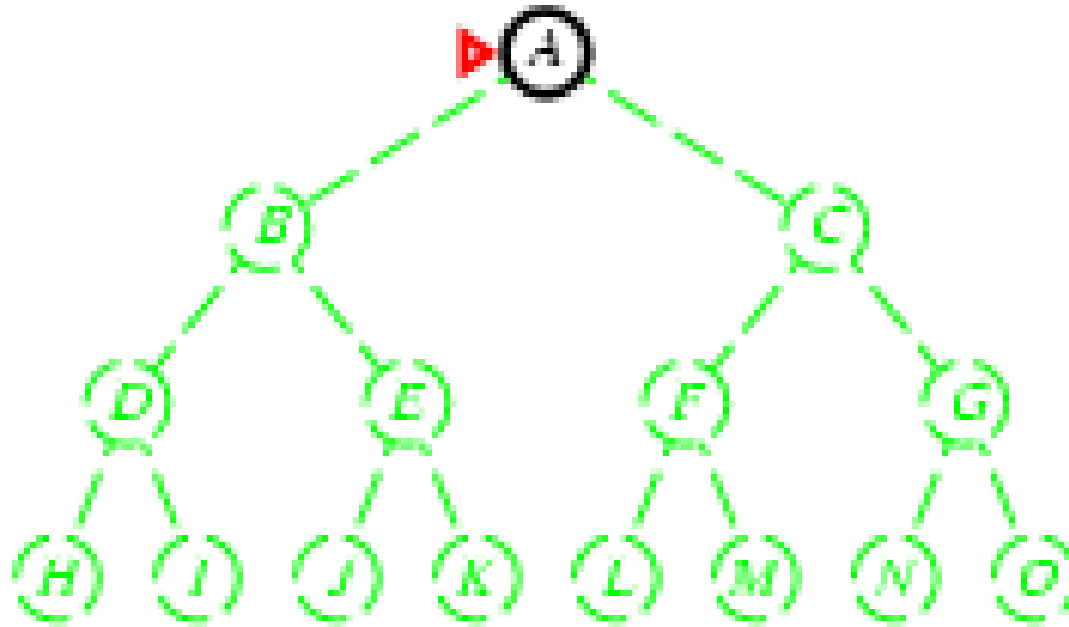
# Properties of breadth-first search

---

- Complete? Yes (if  $b$  is finite)
- Time?  $1+b+b^2+b^3+\dots +b^d + b(b^d-1) = O(b^{d+1})$
- Space?  $O(b^{d+1})$  (keeps every node in memory)
- Optimal? Yes (if cost = 1 per step)
- Space is the bigger problem (more than time)

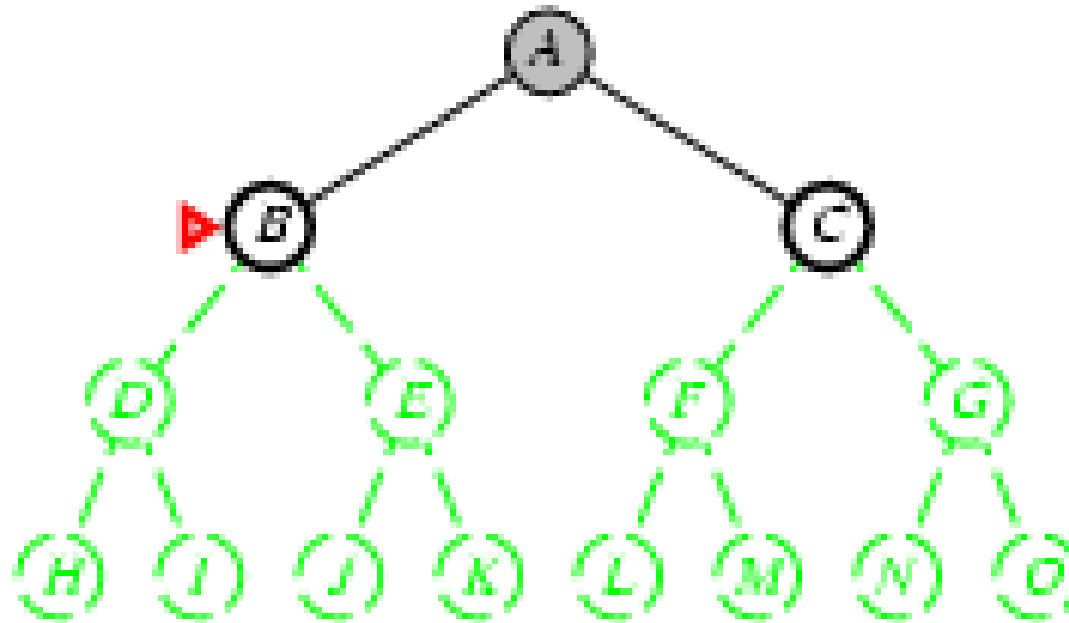
# Depth-first search

- Expand deepest unexpanded node
- **Implementation:**
  - *fringe* = LIFO queue, i.e., put successors at front



# Depth-first search

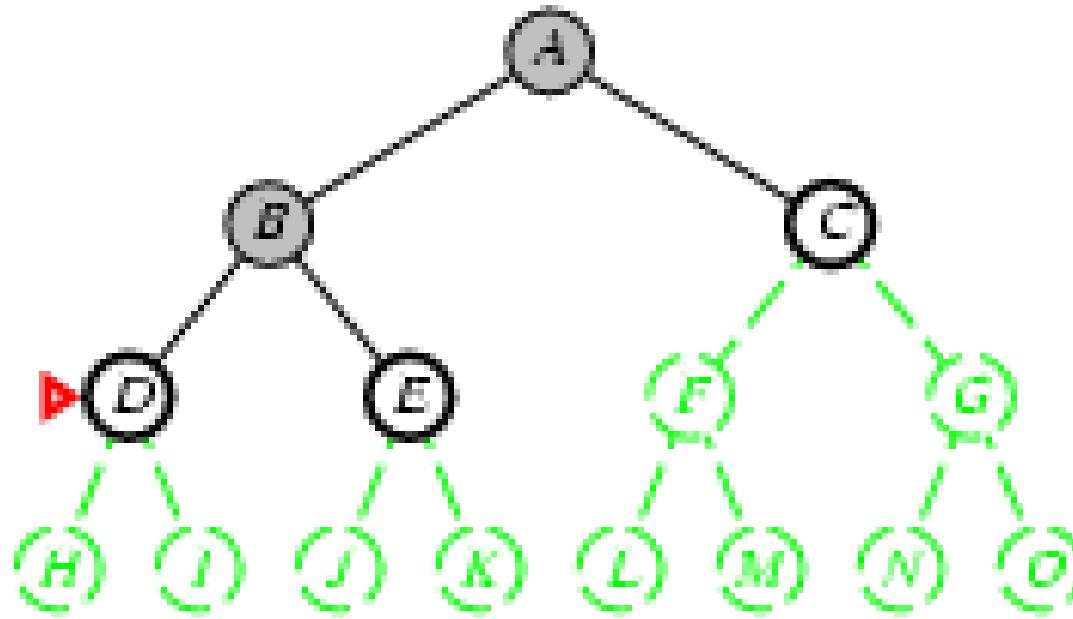
- Expand deepest unexpanded node
- **Implementation:**
  - *fringe* = LIFO queue, i.e., put successors at front





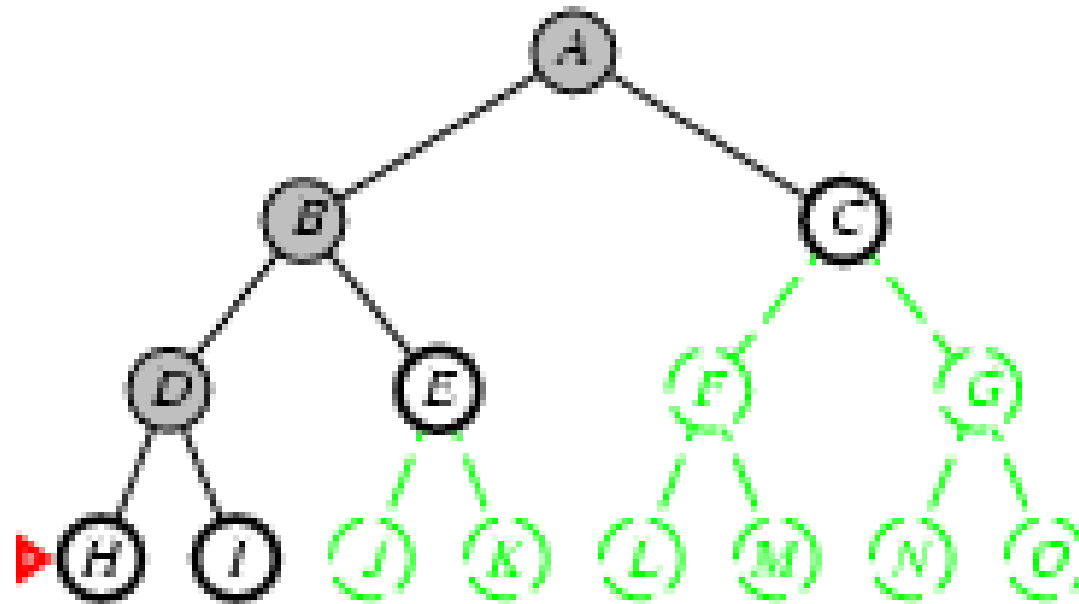
# Depth-first search

- Expand deepest unexpanded node
- **Implementation:**
  - *fringe* = LIFO queue, i.e., put successors at front



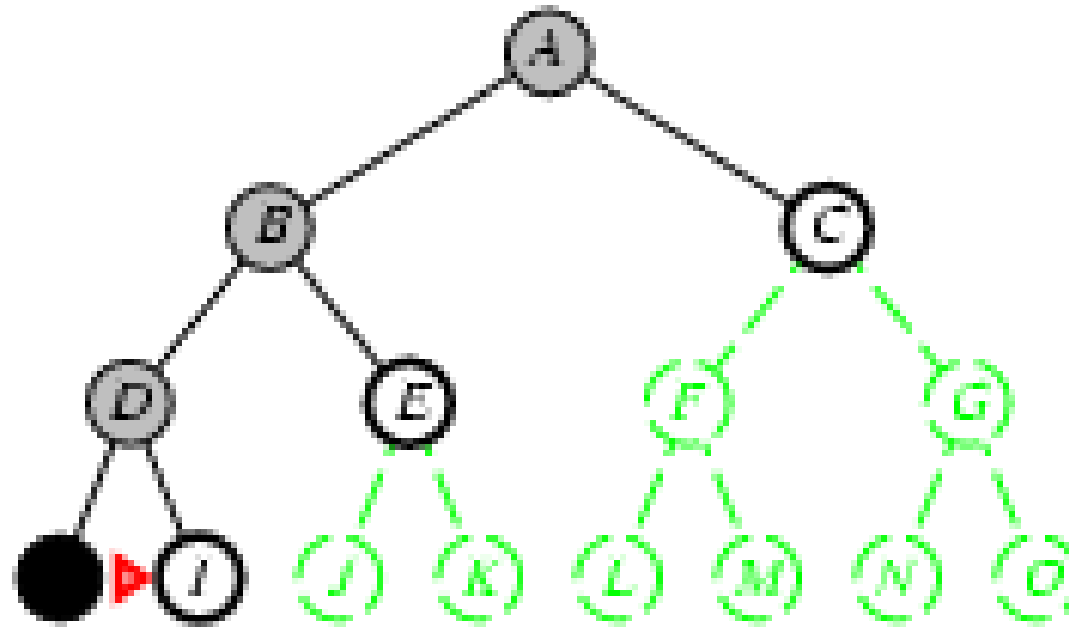
# Depth-first search

- Expand deepest unexpanded node
- **Implementation:**
  - *fringe* = LIFO queue, i.e., put successors at front



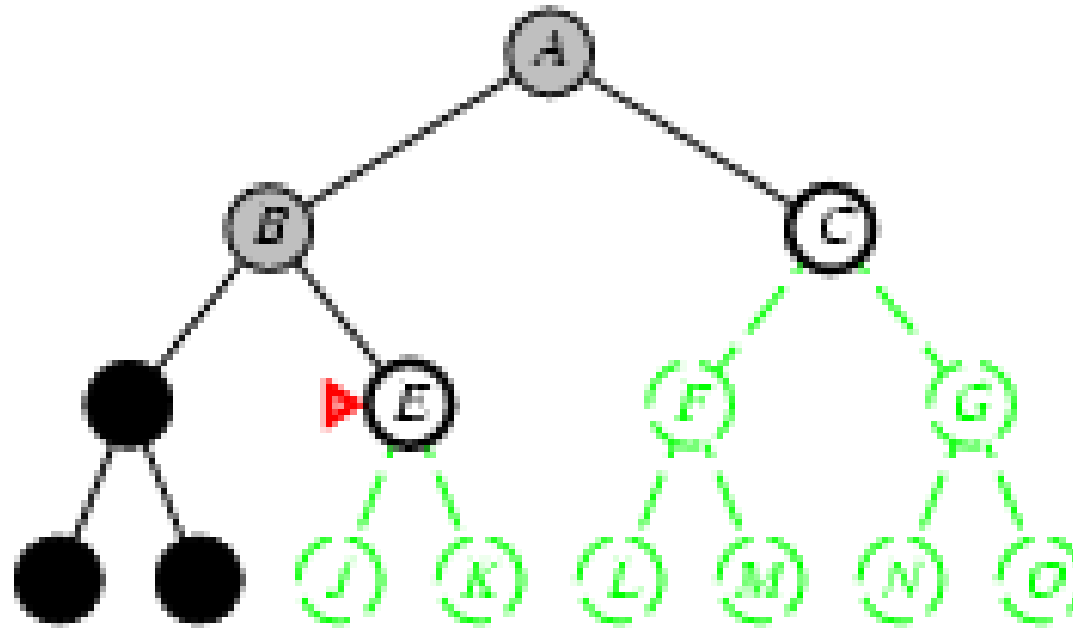
# Depth-first search

- Expand deepest unexpanded node
- **Implementation:**
  - *fringe* = LIFO queue, i.e., put successors at front



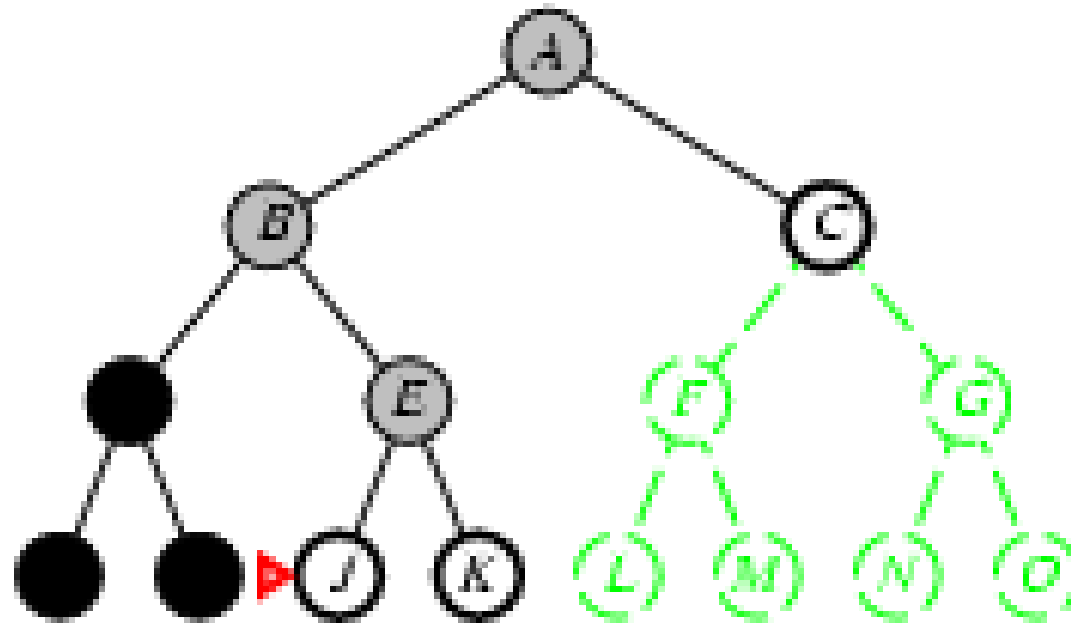
# Depth-first search

- Expand deepest unexpanded node
- **Implementation:**
  - *fringe* = LIFO queue, i.e., put successors at front



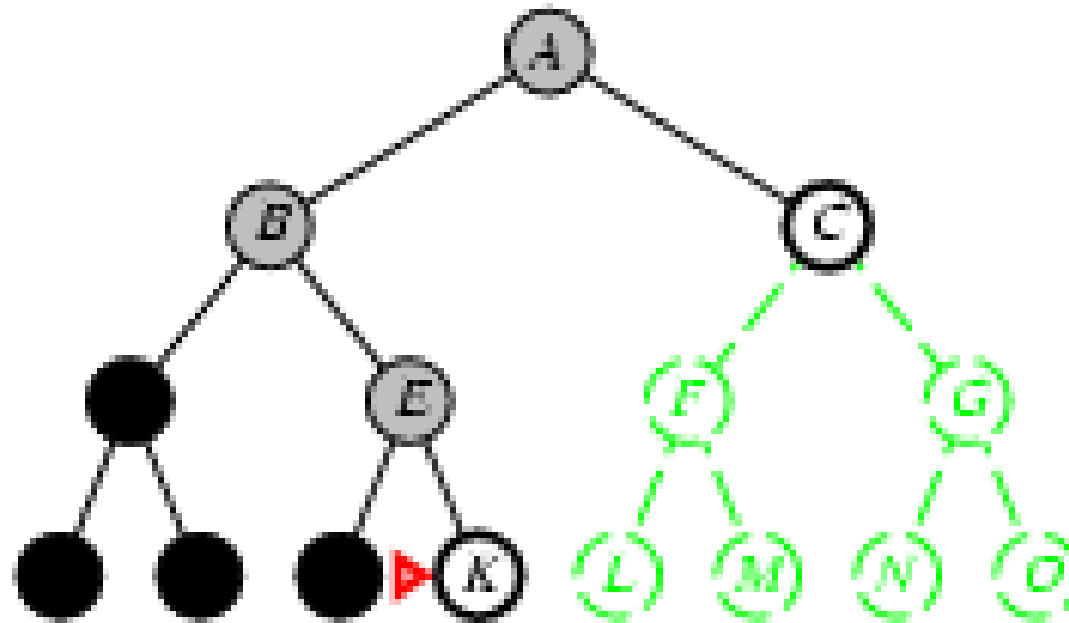
# Depth-first search

- Expand deepest unexpanded node
- **Implementation:**
  - *fringe* = LIFO queue, i.e., put successors at front



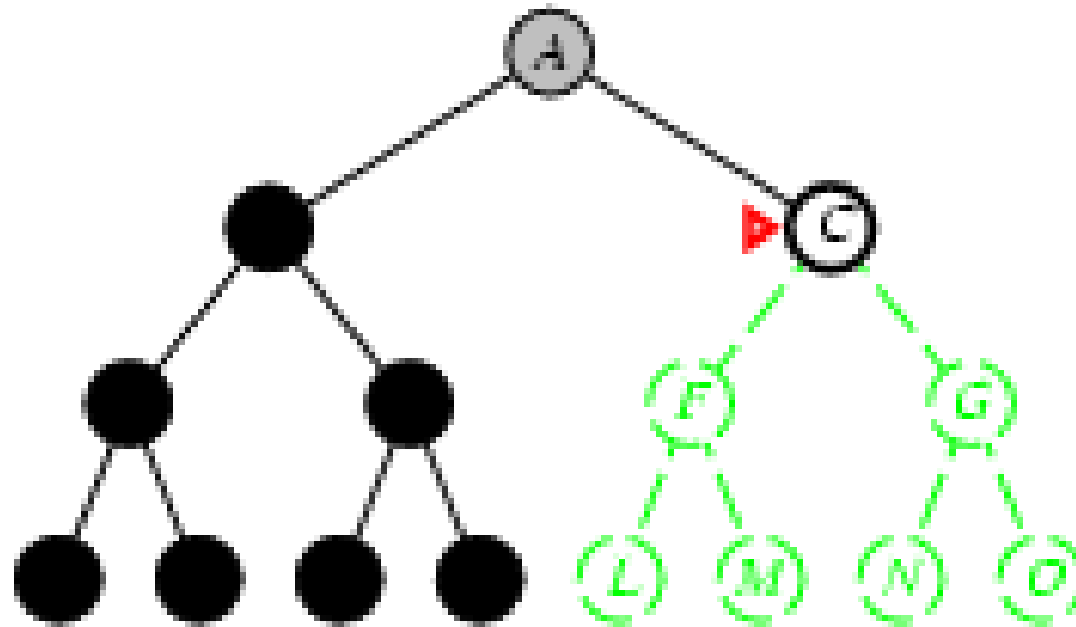
# Depth-first search

- Expand deepest unexpanded node
- **Implementation:**
  - *fringe* = LIFO queue, i.e., put successors at front



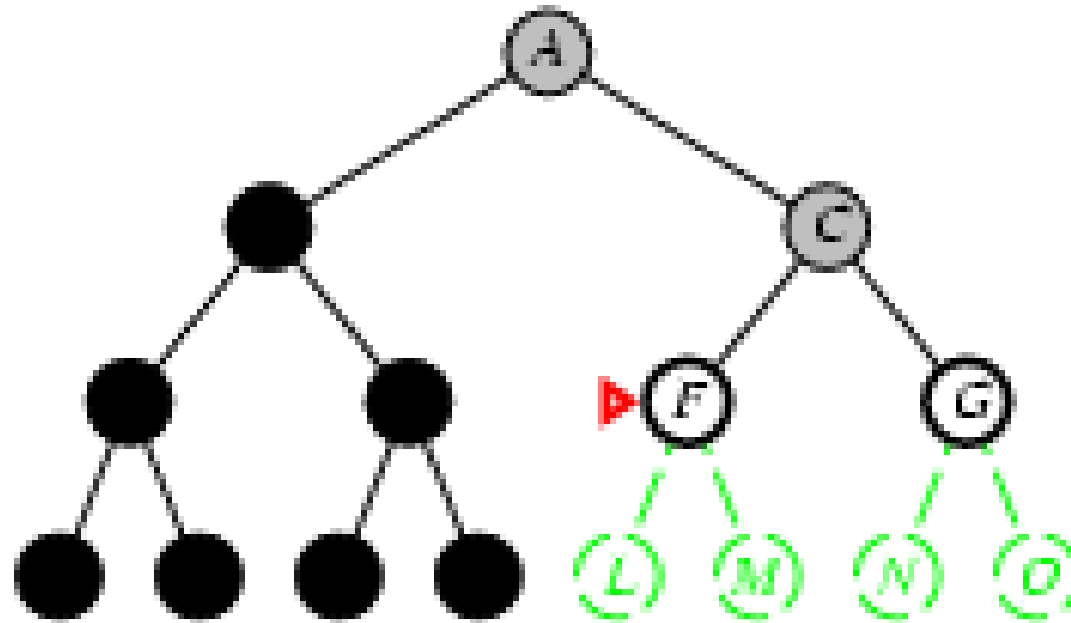
# Depth-first search

- Expand deepest unexpanded node
- **Implementation:**
  - *fringe* = LIFO queue, i.e., put successors at front



# Depth-first search

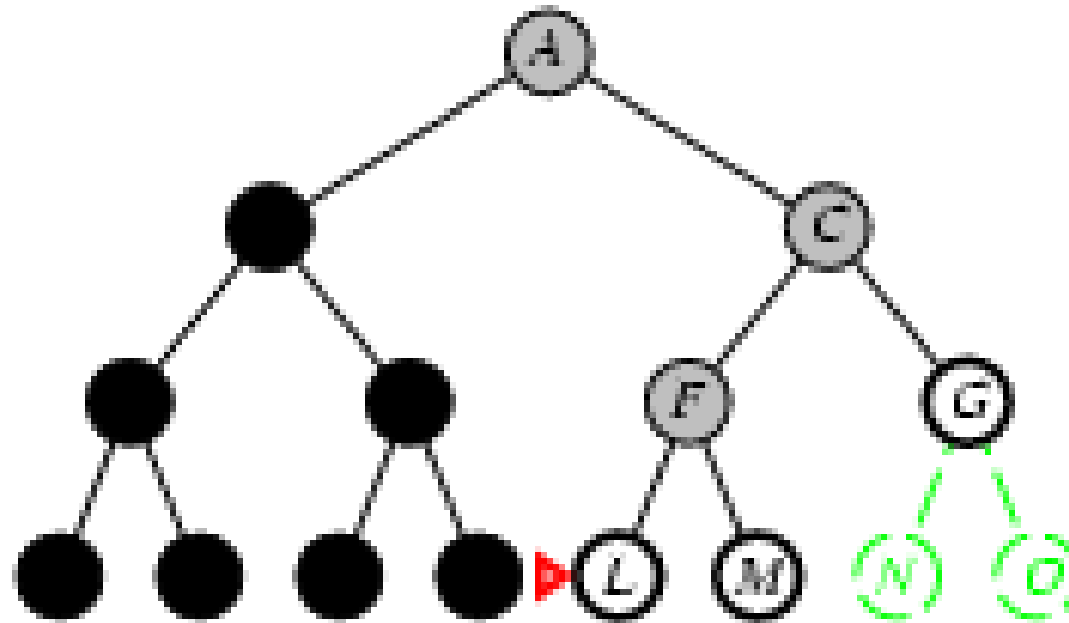
- Expand deepest unexpanded node
- **Implementation:**
  - *fringe* = LIFO queue, i.e., put successors at front





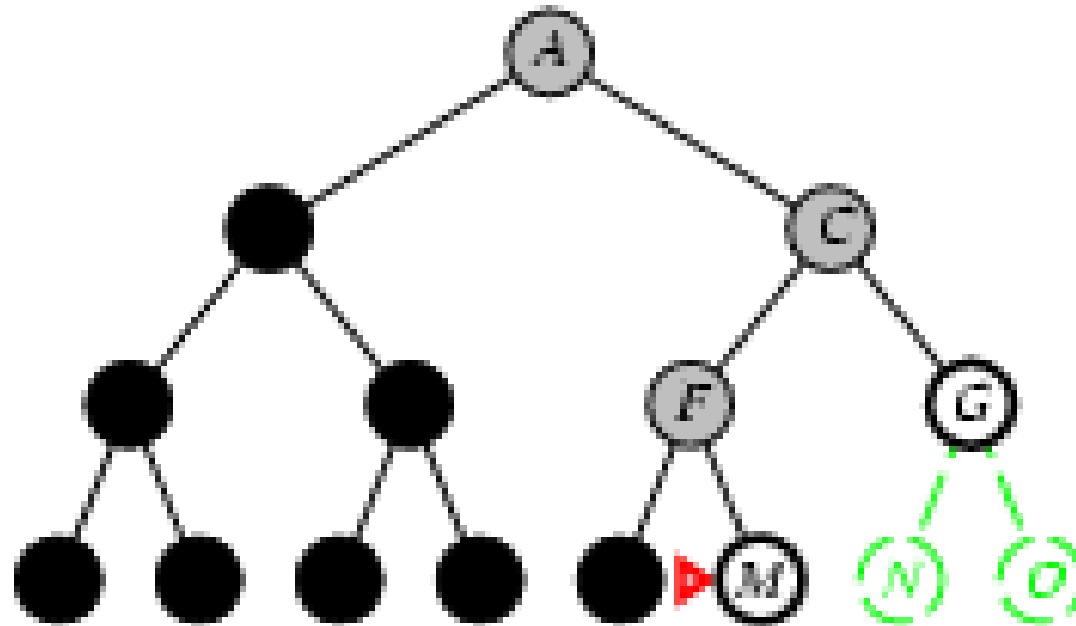
# Depth-first search

- Expand deepest unexpanded node
- **Implementation:**
  - *fringe* = LIFO queue, i.e., put successors at front
  -



# Depth-first search

- Expand deepest unexpanded node
- **Implementation:**
  - *fringe* = LIFO queue, i.e., put successors at front



# Properties of depth-first search

---

- **Complete?** No: fails in **infinite-depth** spaces, spaces with loops
  - Modify to avoid repeated states along path
    - $\rightarrow$  **complete in finite spaces**
- **Time?**  $O(b^m)$ : terrible if  $m$  is much larger than  $d$ 
  - but if solutions are dense, may be much faster than breadth-first
- **Space?**  $O(bm)$ , i.e., **linear space!**
- **Optimal?** No

# Depth-limited search

---

= depth-first search with depth limit  $l$ ,  
i.e., nodes at depth  $l$  have no successors

# Iterative deepening search

---

Function Iterative\_Deepening\_Search(*problem*) return *solution* or *failure*

Inputs: *problem*, a problem

For *depth*  $\leftarrow 0$  to  $\infty$  do

*result*  $\leftarrow$  Depth\_Limited\_Search (*problem*, *depth*)

    if *result*  $\neq$  cutoff then return *result*

# Iterative deepening search $l = 0$

---

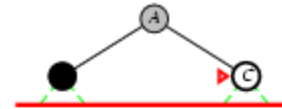
Limit = 0



# Iterative deepening search $l = 1$

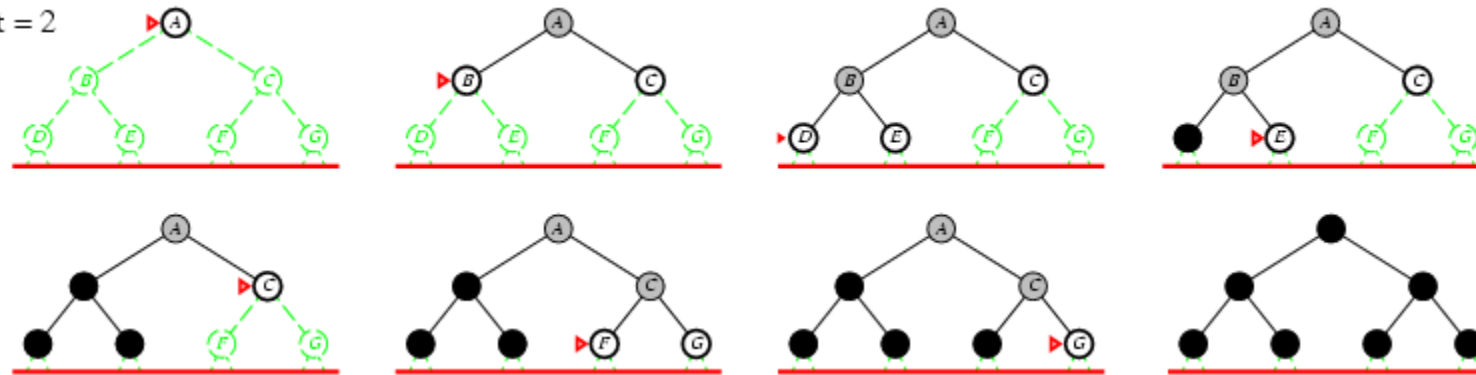
---

Limit = 1



# Iterative deepening search $l=2$

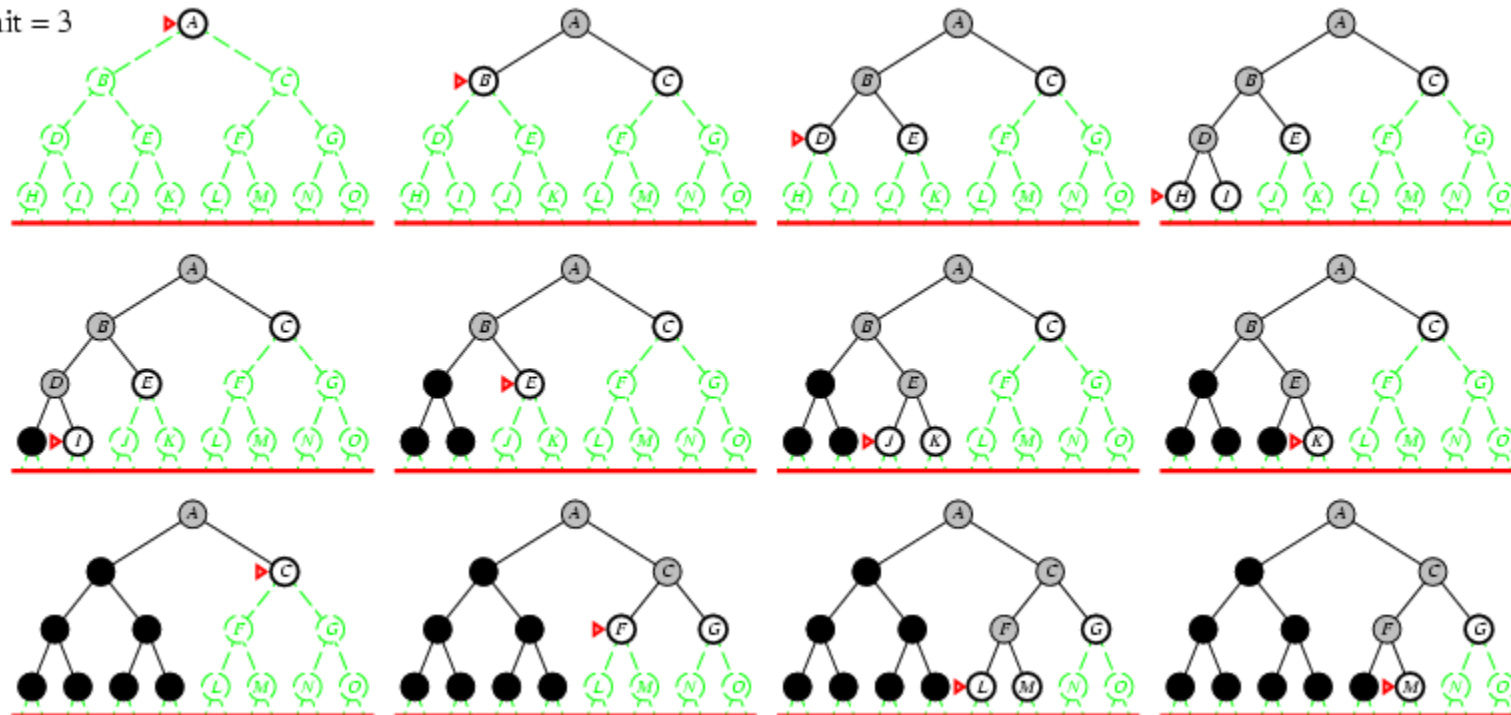
Limit = 2





# Iterative deepening search $l=3$

Limit = 3



# Iterative deepening search

---

- Number of nodes generated in a **depth-limited search** to depth  $d$  with branching factor  $b$ :
  - $N_{DLS} = b^0 + b^1 + b^2 + \dots + b^{d-2} + b^{d-1} + b^d$
- Number of nodes generated in **an iterative deepening search** to depth  $d$  with branching factor  $b$ :
  - $N_{IDS} = (d+1)b^0 + d b^1 + (d-1)b^2 + \dots + 3b^{d-2} + 2b^{d-1} + 1b^d$
- For  $b = 10, d = 5$ ,
  - $N_{DLS} = 1 + 10 + 100 + 1,000 + 10,000 + 100,000 = 111,111$
  - $N_{IDS} = 6 + 50 + 400 + 3,000 + 20,000 + 100,000 = 123,456$
- Overhead =  $(123,456 - 111,111)/111,111 = 11\%$

# Properties of iterative deepening search

---

- Complete? Yes
- Time?  $(d+1)b^0 + d b^1 + (d-1)b^2 + \dots + b^d = O(b^d)$
- Space?  $O(bd)$
- Optimal? Yes, if step cost = 1

# Summary of blind search algorithms

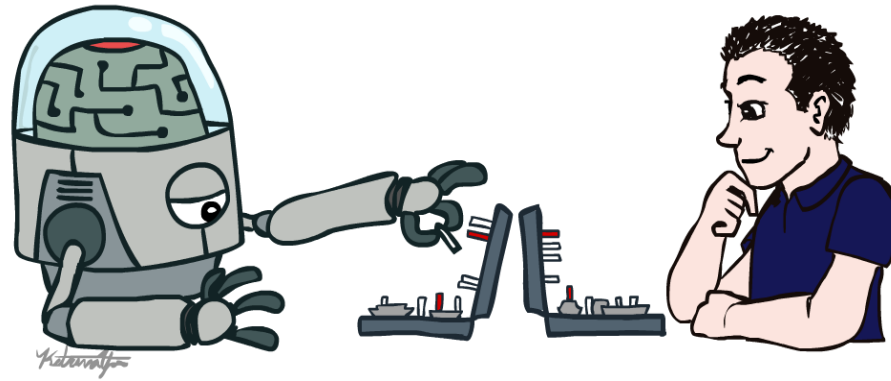
---

Criterion	Breadth First	Depth First	Depth Limited	Iterative Deepening
Complete?	Yes	No	No	Yes
Time	$O(b^{d+1})$	$O(b^m)$	$O(b^l)$	$O(b^d)$
Space	$O(b^{d+1})$	$O(bm)$	$O(bl)$	$O(bd)$
Optimal	Yes	No	No	Yes

# Next class?

---

- Informed search
- **Heuristic search**



Thanks!