

Solving Cryptarithmic Puzzles by Logic Programming

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Abstract—As a personal interest of study, I tried a logic programming approach towards the problem solving of cryptarithmic puzzles that are commonly discussed as a subcategory of constraint satisfaction problems in the literature of artificial intelligence. While there are possibly several methods capable of solving constraint satisfaction problems, I took into consideration the efficiency as well as the completeness that will identify all possible solutions under the specified constraints and exclude trivial and useless solutions from the perspective of real-life practice. In this paper, I demonstrated an approach that can be adapted to solve most of the constraint satisfaction problems especially within the context of cryptarithmic puzzles. This method will also perform forward checking to have early backtracking and prevent searching the entire search tree exhaustively.

Index Terms—Cryptarithmic puzzle, Constraint Satisfaction Problem, Forward Checking, Early Backtracking

I. INTRODUCTION

In the literature of Artificial Intelligence, cryptarithmic puzzles are generally discussed as a kind of the Constraint Satisfaction Problems (CSPs) in which a solution to a given problem is represented by a problem state that meets of all the problem constraints. In the context of presenting a cryptarithmic puzzle, the numerical values involved in an arithmetic computation are encrypted and represented not by numerical numbers, formed by digits from 0 to 9, but by encrypted alphabetical letters.

By concerning the real-life practice of data encryption and mathematical correctness, the solutions to a given puzzle have to comply with the following constraints:

- Each letter involved in the computation is representing a digit.
- No two letters are representing the same digit.
- After replacing each letter by its corresponding digit, the resultant value is mathematically correct.

As an instance of cryptographical problems, let's consider the following example:

```

  WAN
+ LAN
-----
 BOB

```

A possible solution to this problem is replacing A by 7, replacing B by 8, replacing L by 3, replacing N by 9, replacing O by 5, and replacing W by 4. As a result, this solution is viewed and verified as:

A=7, B=8, L=3, N=9, O=5, W=4

```

  WAN      479
+ LAN      + 379
-----
  BOB      858

```

In this paper, I applied a logical programming approach to search for all possible solutions to the above instance of cryptarithmic puzzles. While most of the constraint satisfaction problems can be solved by brute force methods, such as generate and test, this inefficient approach is usually not preferred in the society artificial intelligence. To prevent exhaustively searching the entire search tree, a forward checking method is incorporated to have early backtracking.

It is also worthy of mentioning that the problem constraints are purposefully specified to exclude the trivial solution that is simply replacing all letters by 0's such as:

```

  000
+ 000
-----
  000

```

From the standing point of cryptography, this trivial encryption and decryption are useless and not of much practical interest in real-life message security.

II. SOME PRELIMINARY ANALYSIS OF POSSIBLE METHODOLOGIES

Based on the general approaches of solving constraint satisfaction problems, the following analysis discussed two possible methods that are applicable of solving the aforementioned instance of puzzle.

A. The Brute Force Methodology

Let's start with solving a cryptarithmic puzzle by using the brute force methods. This is by far the most intuitive but not so intelligent method. While it is not a heuristically informed method and often criticized by its inefficiency, the generate and test method is a very fundamental brute force method that is capable of solving problems in which solutions can be found within a given problem's search space. With no exception, this method is capable enough of solving any cryptarithmic puzzle without concerning its inefficiency. The very core of a brute force problem-solving method is exhaustively searching through the entire search space that consists of all possible

problem states and screen out those states that are not meeting all problem constraints.

Based on the instance that has been illustrated in this paper, the entire search space can be represented by the following search tree that is partially represented in Figure 1. Within this search tree, every path from the root to a leaf node is representing a possible problem state that consists of possible encryption consists of replacing every letter by the digit suggested along the path. Also, for the purpose of following the mathematical convention that starts the computation of addition from the 1's place and then goes on to the 10's place, 100's place, and so on. This search tree starts form possible encryption of the root N, then B, A, O, W, and finally ends at L.

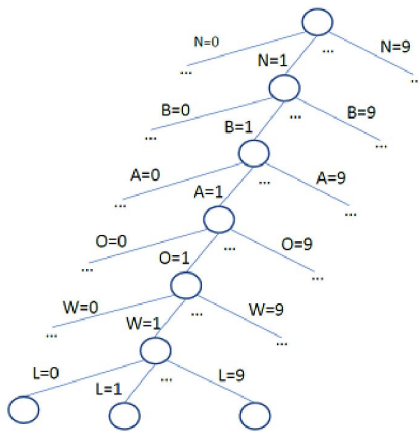


Fig. 1. Part of the Search Tree

The complete search tree comes with a height of 6 and a branching factor of 10. So, the entire search space has 1,000,000 paths. For the purpose of identifying all possible solutions, this brute force method can verify the correctness of each path by replacing letters with digits that are suggested along the path and screen out paths that are not meeting all the problem constraints.

Although this may appear to be a valid approach to solve the problem, it is simply not efficient and involves the examination of unnecessary replacements of letters by digits.

B. Applying Forward Checking Methods

While the brute force search is valid and will not miss any possible solution, it does involve many unnecessary overheads on the way of examining each path. This overhead can be avoided by applying the idea of forward checking to have early backtracking before processing the entire path. As soon as any of the letter replacement is found to be violating any problem constraint the rest of the path can be abandoned without any further examination. As an illustration of this forward checking, let's consider the path consists of $N=1$, $B=1$, $A=2$, $O=3$, $W=4$, and $L=5$. Right after seeing that B is repeating the digit 1, we notice the violation of "no two letters are representing the same digit." So, the rest of the path can

be skipped from any further examination and the search will go on to the verification of the next path

A comprehensive verification of the conformance with problem constraints can be described in the following sequence:

By examining the computation of 1's place, we know that:

- N can be replaced by any digit.
- B can be replaced by a digit that is not used to replace N.
- The remainder of $(N + N) \div 10$ must be equal to B.

By examining the computation from 1's place to 10's place, we know that:

- N can be replaced by any digit.
- B can be replaced by a digit that is not used to replace N.
- A can be replaced by a digit that is not used to replace N and B.
- O can be replaced by a digit that is not used to replace N, B and A.
- The remainder of $(N + N + (A \times 10) + (A \times 10)) \div 100$ must be equal to $(B + (O \times 10))$.

By examining the computation from 1's place to 100's place, we know that:

- N can be replaced by any digit.
- B can be replaced by a digit that is not used to replace N.
- A can be replaced by a digit that is not used to replace N and B.
- O can be replaced by a digit that is not used to replace N, B and A.
- W can be replaced by a digit that is not used to replace N, B, A and O.
- L can be replaced by a digit that is not used to replace N, B, A, O and W.
- Since there is no carrying from the 100's place to the 1000's place, $(N + N + (A \times 10) + (A \times 10) + (W \times 100) + (L \times 100))$ must be equal to $(B + (O \times 10) + (B \times 100))$.

III. THE IMPLEMENTATION IN CLIPS PROGRAMMING LANGUAGE

While there are many programming languages that can be adopted to implement the aforementioned problem-solving method, by concerning the advantage of having a built-in inference engine, I focused my selection on languages in the paradigm of logic programming and adopted the CLIPS programming language [1] [2] [3].

The initial fact-base comes with all letters and digits within the problem domain. In CLIPS syntax they are defined as a group of initial facts as follows:

```
(defacts letters-and-digits
  (letter A)
  (letter B)
  (letter L)
  (letter N)
  (letter O)
  (letter W)
  (digit 0)
```

```
(digit 1)
(digit 2)
(digit 3)
(digit 4)
(digit 5)
(digit 6)
(digit 7)
(digit 8)
(digit 9))
```

In addition to the initial contents, all possible replacements of letters by digits are added to the fact-base by the following knowledge rule:

```
(defrule all-replacements
  (letter ?letter)
  (digit ?digit)
  =>
  (assert (replace ?letter ?digit)))
```

The real forward checking is implemented by the following knowledge rule:

```
(defrule forward-checking

  ;checking the 1's place
  (replace N ?n)
  (replace B ?b&~?n)
  (test (= (mod (+ ?n ?n) 10) ?b))

  ;checking up to 10's place
  (replace A ?a&~?n&~?b)
  (replace O ?o&~?n&~?b&~?a)
  (test (= (mod (+ ?n ?n (* ?a 10) (* ?a 10)) 100)
    (+ (* ?o 10) ?b)))

  ;checking up to 100's place
  (replace W ?w&~?n&~?b&~?a&~?o)
  (replace L ?l&~?n&~?b&~?a&~?o&~?w)
  (test (= (+ ?n ?n (* ?a 10) (* ?a 10) (* ?w 100)
    (* ?l 100))
    (+ (* ?b 100) (* ?o 10) ?b)))

  =>

  ;display all solutions
  (printout t "A=" ?a " , B=" ?b " , L=" ?l " , N=" ?n
    " , O=" ?o " , W=" ?w crlf cefl)
  (printout t "      WAN      " ?w ?a ?n
    crlf)
  (printout t "      + LAN      " + " ?l ?a ?n
    crlf)
  (printout t "      -----    ==>    -----" crlf)
  (printout t "      BOB      " ?b ?o ?b
    crlf crlf) )
```

IV. THE LIST OF ALL POSSIBLE SOLUTIONS

All possible solutions to the given cryptarithmic problem are listed as follows:

A=7, B=8, L=3, N=9, O=5, W=4

```
      WAN      479
+   LAN      + 379
-----    ==>    -----
      BOB      858
```

A=7, B=8, L=4, N=9, O=5, W=3

```
      WAN      379
+   LAN      + 479
-----    ==>    -----
      BOB      858
```

A=3, B=6, L=2, N=8, O=7, W=4

```
      WAN      438
+   LAN      + 238
-----    ==>    -----
      BOB      676
```

A=3, B=8, L=2, N=9, O=7, W=6

```
      WAN      639
+   LAN      + 239
-----    ==>    -----
      BOB      878
```

A=6, B=8, L=2, N=9, O=3, W=5

```
      WAN      569
+   LAN      + 269
-----    ==>    -----
      BOB      838
```

A=7, B=6, L=2, N=8, O=5, W=3

```
      WAN      378
+   LAN      + 278
-----    ==>    -----
      BOB      656
```

A=7, B=6, L=3, N=8, O=5, W=2

```
      WAN      278
+   LAN      + 378
-----    ==>    -----
      BOB      656
```

A=3, B=6, L=4, N=8, O=7, W=2

```
      WAN      238
+   LAN      + 438
-----    ==>    -----
      BOB      676
```

A=3, B=8, L=6, N=9, O=7, W=2

```
      WAN      239
+   LAN      + 639
-----    ==>    -----
      BOB      878
```

A=6, B=8, L=5, N=9, O=3, W=2

```
      WAN      269
+   LAN      + 569
-----    ==>    -----
      BOB      838
```

A=1, B=8, L=6, N=9, O=3, W=2

```
      WAN      219
+   LAN      + 619
-----    ==>    -----
      BOB      838
```

A=1, B=8, L=2, N=9, O=3, W=6

```
      WAN      619
+   LAN      + 219
-----    ==>    -----
      BOB      838
```

A=1, B=6, L=4, N=8, O=3, W=2

```
      WAN      218
```

+ LAN		+ 418
-----	====>	-----
BOB		636

A=1, B=6, L=2, N=8, O=3, W=4

WAN		418
+ LAN		+ 218
-----	====>	-----
BOB		636

A=1, B=8, L=5, N=4, O=2, W=3

WAN		314
+ LAN		+ 514
-----	====>	-----
BOB		828

A=1, B=8, L=3, N=4, O=2, W=5

WAN		514
+ LAN		+ 314
-----	====>	-----
BOB		828

A=7, B=8, L=1, N=9, O=5, W=6

WAN		679
+ LAN		+ 179
-----	====>	-----
BOB		858

A=4, B=6, L=1, N=8, O=9, W=5

WAN		548
+ LAN		+ 148
-----	====>	-----
BOB		696

A=7, B=6, L=1, N=8, O=5, W=4

WAN		478
+ LAN		+ 178
-----	====>	-----
BOB		656

A=3, B=8, L=1, N=4, O=6, W=7

WAN		734
+ LAN		+ 134
-----	====>	-----
BOB		868

A=3, B=6, L=1, N=8, O=7, W=5

WAN		538
+ LAN		+ 138
-----	====>	-----
BOB		676

A=9, B=6, L=1, N=3, O=8, W=4

WAN		493
+ LAN		+ 193
-----	====>	-----
BOB		686

A=4, B=6, L=1, N=3, O=8, W=5

WAN		543
+ LAN		+ 143
-----	====>	-----
BOB		686

A=2, B=6, L=1, N=3, O=4, W=5

WAN		523
+ LAN		+ 123
-----	====>	-----
BOB		646

A=2, B=4, L=1, N=7, O=5, W=3

WAN		327
+ LAN		+ 127
-----	====>	-----
BOB		454

A=2, B=8, L=1, N=9, O=5, W=7

WAN		729
+ LAN		+ 129
-----	====>	-----
BOB		858

A=6, B=4, L=1, N=7, O=3, W=2

WAN		267
+ LAN		+ 167
-----	====>	-----
BOB		434

A=5, B=6, L=3, N=8, O=1, W=2

WAN		258
+ LAN		+ 358
-----	====>	-----
BOB		616

A=5, B=6, L=2, N=8, O=1, W=3

WAN		358
+ LAN		+ 258
-----	====>	-----
BOB		616

A=5, B=8, L=4, N=9, O=1, W=3

WAN		359
+ LAN		+ 459
-----	====>	-----
BOB		818

A=5, B=8, L=3, N=9, O=1, W=4

WAN		459
+ LAN		+ 359
-----	====>	-----
BOB		818

A=7, B=6, L=4, N=8, O=5, W=1

WAN		178
+ LAN		+ 478
-----	====>	-----
BOB		656

A=7, B=8, L=6, N=9, O=5, W=1

WAN		179
+ LAN		+ 679
-----	====>	-----
BOB		858

A=4, B=6, L=5, N=8, O=9, W=1

WAN		148
+ LAN		+ 548

```

-----  ===>  -----
      BOB                696
A=3, B=8, L=7, N=4, O=6, W=1

      WAN                134
+ LAN                + 734
-----  ===>  -----
      BOB                868
A=3, B=6, L=5, N=8, O=7, W=1

      WAN                138
+ LAN                + 538
-----  ===>  -----
      BOB                676
A=9, B=6, L=4, N=3, O=8, W=1

      WAN                193
+ LAN                + 493
-----  ===>  -----
      BOB                686
A=4, B=6, L=5, N=3, O=8, W=1

      WAN                143
+ LAN                + 543
-----  ===>  -----
      BOB                686
A=6, B=4, L=2, N=7, O=3, W=1

      WAN                167
+ LAN                + 267
-----  ===>  -----
      BOB                434
A=2, B=6, L=5, N=3, O=4, W=1

      WAN                123
+ LAN                + 523
-----  ===>  -----
      BOB                646
A=2, B=4, L=3, N=7, O=5, W=1

      WAN                127
+ LAN                + 327
-----  ===>  -----
      BOB                454
A=2, B=8, L=7, N=9, O=5, W=1

      WAN                129
+ LAN                + 729
-----  ===>  -----
      BOB                858
A=0, B=8, L=6, N=9, O=1, W=2

      WAN                209
+ LAN                + 609
-----  ===>  -----
      BOB                818
A=0, B=8, L=5, N=9, O=1, W=3

      WAN                309
+ LAN                + 509
-----  ===>  -----
      BOB                818
A=0, B=8, L=3, N=9, O=1, W=5

```

```

      WAN                509
+ LAN                + 309
-----  ===>  -----
      BOB                818
A=0, B=8, L=2, N=9, O=1, W=6

      WAN                609
+ LAN                + 209
-----  ===>  -----
      BOB                818
A=0, B=6, L=4, N=8, O=1, W=2

      WAN                208
+ LAN                + 408
-----  ===>  -----
      BOB                616
A=0, B=6, L=2, N=8, O=1, W=4

      WAN                408
+ LAN                + 208
-----  ===>  -----
      BOB                616
A=9, B=6, L=0, N=3, O=8, W=5

      WAN                593
+ LAN                + 093
-----  ===>  -----
      BOB                686
A=7, B=6, L=0, N=3, O=4, W=5

      WAN                573
+ LAN                + 073
-----  ===>  -----
      BOB                646
A=6, B=8, L=0, N=9, O=3, W=7

      WAN                769
+ LAN                + 069
-----  ===>  -----
      BOB                838
A=9, B=4, L=0, N=2, O=8, W=3

      WAN                392
+ LAN                + 092
-----  ===>  -----
      BOB                484
A=8, B=4, L=0, N=2, O=6, W=3

      WAN                382
+ LAN                + 082
-----  ===>  -----
      BOB                464
A=6, B=8, L=0, N=4, O=2, W=7

      WAN                764
+ LAN                + 064
-----  ===>  -----
      BOB                828
A=5, B=4, L=0, N=7, O=1, W=3

      WAN                357
+ LAN                + 057
-----  ===>  -----

```

BOB	414
-----	-----

A=5, B=8, L=0, N=9, O=1, W=7

WAN	759
+ LAN	+ 059
-----	=====
BOB	818

A=7, B=2, L=0, N=6, O=5, W=1

WAN	176
+ LAN	+ 076
-----	=====
BOB	252

A=8, B=2, L=0, N=6, O=7, W=1

WAN	186
+ LAN	+ 086
-----	=====
BOB	272

A=5, B=8, L=6, N=4, O=0, W=1

WAN	154
+ LAN	+ 654
-----	=====
BOB	808

A=5, B=8, L=1, N=4, O=0, W=6

WAN	654
+ LAN	+ 154
-----	=====
BOB	808

A=5, B=6, L=4, N=3, O=0, W=1

WAN	153
+ LAN	+ 453
-----	=====
BOB	606

A=5, B=6, L=1, N=3, O=0, W=4

WAN	453
+ LAN	+ 153
-----	=====
BOB	606

A=9, B=6, L=5, N=3, O=8, W=0

WAN	093
+ LAN	+ 593
-----	=====
BOB	686

A=7, B=6, L=5, N=3, O=4, W=0

WAN	073
+ LAN	+ 573
-----	=====
BOB	646

A=6, B=8, L=7, N=9, O=3, W=0

WAN	069
+ LAN	+ 769
-----	=====
BOB	838

A=8, B=2, L=1, N=6, O=7, W=0

WAN	086
+ LAN	+ 186
-----	=====
BOB	272

A=7, B=2, L=1, N=6, O=5, W=0

WAN	076
+ LAN	+ 176
-----	=====
BOB	252

A=9, B=4, L=3, N=2, O=8, W=0

WAN	092
+ LAN	+ 392
-----	=====
BOB	484

A=8, B=4, L=3, N=2, O=6, W=0

WAN	082
+ LAN	+ 382
-----	=====
BOB	464

A=6, B=8, L=7, N=4, O=2, W=0

WAN	064
+ LAN	+ 764
-----	=====
BOB	828

A=5, B=4, L=3, N=7, O=1, W=0

WAN	057
+ LAN	+ 357
-----	=====
BOB	414

A=5, B=8, L=7, N=9, O=1, W=0

WAN	059
+ LAN	+ 759
-----	=====
BOB	818

V. SUMMARY

Constraint satisfaction problems have been challenging problem-solvers for a long time for the sake of mental exercises. This paper demonstrated a logical programming approach that can be generalized to solve all cryptarithmic puzzles. While this paper is not creating any new theory or methodology, it can be adopted as case study in artificial intelligence-related.

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