

Assignment 11

Task1

Explain the below concepts with an example in brief.

- Nosql Databases
- Types of Nosql Databases
- CAP Theorem
- HBase Architecture
- HBase vs RDBMS

Nosql Databases

- NoSQL stands for Not only SQL. NoSQL are non-relational, open source, schema less and distributed database.
- They are highly popular today because of their ability to scale out or scale horizontally and flexibility to deal with a variety of data: structured, semi-structured and unstructured data.

• Types of Nosql Databases

There are 4 types of NoSQL :

1. Key-Value pair /Key-Value store:

It maintains data in key-value pairs. It uses a hash table in which there exists a unique key and a pointer to a particular item of data.

e.g.

KeyValueFirst NameSachinLast NameGoradeAddressMumbai, India

Amazon DynamoDB, LinkedIn uses key-value store NoSQL database.

2. Document store:

It maintains data in collections constituted of documents.

It pair each key with a complex data structure known as a document.

Documents can contain many different key-value pairs, or key-array pairs, or even nested documents.

```
e.g.
{
    _id: ObjectId(fdf789)
    title: 'MongoDB Overview',
    description: 'MongoDB is no sql database',
    tags: ['mongodb', 'database', 'NoSQL'],
}
```

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Database: MongoDB, Apache CouchDB, Couchbase uses Document store.

3. Column:

Each storage block has data from only one column. It stores data using a column oriented model.

e.g. Databases: Cassandra, HBase

Facebook, eBay uses Column NoSQL.

4. Graph:

A Graph stores data in nodes. They are also called as Network database.

e.g. Databases: Neo4j, hyperGraphDB.

Walmart uses Graph NoSQL.

• CAP Theorem

The CAP theorem applies to distributed systems that store state. It states that networked shared-data systems can only guarantee or strongly support two of the following three properties:

Consistency -

A guarantee that every node in a distributed cluster returns the same, most recent, successful write. Consistency refers to every client having the same view of the data.

Availability –

Every non-failing node returns a response for all read and write requests in a reasonable amount of time. To be available, every node on (either side of a network partition) must be able to respond in a reasonable amount of time.

Partition Tolerant –

The system continues to function and upholds its consistency guarantees in spite of network partitions. Network partitions are a fact of life. Distributed systems guaranteeing partition tolerance can gracefully recover from partitions once the partition heals.

The CAP theorem categorizes systems into three categories. We can have system with one of these three categories:

• CP (Consistent and Partition Tolerant) -

CP is referring to a category of systems where availability is sacrificed only in the case of a network partition.

• CA (Consistent and Available) -

CA systems are consistent and available systems in the absence of any network partition.

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• AP (Available and Partition Tolerant) -

These are systems that are available and partition tolerant but cannot guarantee consistency.

In case of a network partition (a rare occurrence) one needs to choose between Availability and partition tolerance. In any networked shared-data systems partition tolerance is a must. Network partitions and dropped messages are a fact of life and must be handled appropriately. Consequently, system designers must choose between consistency and availability. A network partition forces designers to either choose perfect consistency or perfect availability. Picking consistency means not being able to answer a client's query as the system cannot guarantee to return the most recent write. This sacrifices availability. Network partition forces nonfailing nodes to reject clients' requests as these nodes cannot guarantee consistent data. At the opposite end of the spectrum, being available means being able to respond to a client's request but the system cannot guarantee consistency, i.e., the most recent value written. Available systems provide the best possible answer under the given circumstance.

HBase Architecture

As we know, HBase is a column-oriented NoSQL database.

In a column-oriented databases, all the column values are stored together like first column values will be stored together, then the second column values will be stored together and data in other columns are stored in a similar manner.

- When the amount of data is very huge, like in terms of petabytes or exabytes, we use column-oriented approach, because the data of a single column is stored together and can be accessed faster.
- While row-oriented approach comparatively handles less number of rows and columns efficiently, as row-oriented database stores data is a structured format.
- When we need to process and analyse a large set of semi-structured or unstructured data, we use column oriented approach. Such as applications dealing with Online Analytical Processing like data mining, data warehousing, applications including analytics, etc.
- Whereas, Online Transactional Processing such as banking and finance domains which handle structured data and require transactional properties (ACID properties) use row-oriented approach.

HBase Architecture: Components of HBase Architecture

HBase has three major components i.e., **HMaster Server**, **HBase Region Server**, **Regions** and **Zookeeper**.

HBase Architecture: Region

- A table can be divided into a number of regions. A Region is a sorted range of rows storing data between a start key and an end key.
- A Region has a default size of 256MB which can be configured according to the need.
- A Group of regions is served to the clients by a Region Server.

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• A Region Server can serve approximately 1000 regions to the client.

HBase Architecture: HMaster

As in the below image, you can see the HMaster handles a collection of Region Server which resides on DataNode.

- HBase HMaster performs DDL operations (create and delete tables) and assigns regions to the Region servers as you can see in the above image.
- It coordinates and manages the Region Server (similar as NameNode manages DataNode in HDFS).
- It assigns regions to the Region Servers on startup and re-assigns regions to Region Servers during recovery and load balancing.
- It monitors all the Region Server's instances in the cluster (with the help of Zookeeper) and performs recovery activities whenever any Region Server is down.
- It provides an interface for creating, deleting and updating tables.

ZooKeeper – The Coordinator

- Zookeeper acts like a coordinator inside HBase distributed environment. It helps in maintaining server state inside the cluster by communicating through sessions.
- Every Region Server along with HMaster Server sends continuous heartbeat at regular interval to Zookeeper and it checks which server is alive. It also provides server failure notifications so that, recovery measures can be executed.
- The inactive server, which acts as a backup for active server. If the active server fails, it comes for the rescue
- The active HMaster sends heartbeats to the Zookeeper while the inactive HMaster listens for the notification send by active HMaster. If the active HMaster fails to send a heartbeat the session is deleted and the inactive HMaster becomes active.

• HBase vs RDBMS

HBase	RDBMS
1. Column-oriented	1. Row-oriented
2. Flexible schema, add columns on the Fly	2. Fixed schema
3. Does not supports SQL	3. It supports SQL
4. Good for semi-structured data as well as structured	4. Good for structured data only.
data.	
5. Supports Horizontal scaling(scale out). So easy to	5. Supports Vertical scaling(scale up). So difficult
scale.	to scale.
6. It doesn't supports referential integrity.	6. It supports referential integrity.
7. Good with sparse tables	7. Not optimized for sparse tables.
8. Tight integration with MapReduce (MR).	8. No integration with MapReduce (MR).
9. Joins are not optimized and are using MR.	9. Joins are optimized and are not using MR.
10. It does not supports ACID properties. Hence it	10. It supports ACID properties. Hence supports
does not supports transactions.	transactions.

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Task2:

Execute blog present in below link

https://acadgild.com/blog/importtsv-data-from-hdfs-into-hbase/

We have started all hadoop daemons by using command: start-all.sh

Then we have verified that all hadoop daemons are started by using **jps** command.

```
[acadgild@localhost ~]$ start-all.sh
This script is Deprecated. Instead use start-dfs.sh and start-yarn.sh
18/08/18 10:48:13 WARN util.NativeCodeLoader: Unable to load native-hadoop library for your platform... using builtin-java classes where applicable
Starting namenodes on [localhost]
localhost: starting namenode, logging to /home/acadgild/install/hadoop/hadoop-2.6.5/logs/hadoop-acadgild-namenode-localhost.localdomain.o
ut
localhost: starting datanode, logging to /home/acadgild/install/hadoop/hadoop-2.6.5/logs/hadoop-acadgild-datanode-localhost.localdomain.o
ut
Starting secondary namenodes [0.0.0.0]
0.0.0: starting secondarynamenode, logging to /home/acadgild/install/hadoop/hadoop-2.6.5/logs/hadoop-acadgild-secondarynamenode-localhos
t.localdomain.out
18/08/18 10:48:58 WARN util.NativeCodeLoader: Unable to load native-hadoop library for your platform... using builtin-java classes where applicable
starting yarn daemons
starting resourcemanager, logging to /home/acadgild/install/hadoop/hadoop-2.6.5/logs/yarn-acadgild-resourcemanager-localhost.localdomain.out
localhost: starting nodemanager, logging to /home/acadgild/install/hadoop/hadoop-2.6.5/logs/yarn-acadgild-nodemanager-localhost.localdomain.out
You have new mail in /var/spool/mail/acadgild
[acadgild@localhost ~]s
[jps]
372 NodeManager
3755 SecondaryNameNode
3768 ResourceManager
4458 Jps
3451 DataNode
3355 NameNode
```

After this, we have started hbase and job history daemons by using commands: start-hbase.sh and /home/acadgild/install/hadoop/hadoop-2.6.5/sbin/mr-jobhistory-daemon.sh start historyserver

Then we have verified that hbase and job history server daemons are started by using jps command.

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```
[acadgild@localhost ~|s start.hbase.sh localhost: starting zookeeper, logging to /home/acadgild/install/hbase/hbase-1.2.6/logs/hbase-acadgild-zookeeper-localhost.localdomain.out tarting measter, logging to /home/acadgild/install/hbase/hbase-1.2.6/logs/hbase-acadgild-master-localhost.localdomain.out facadgild@localhost ~|s jps 3872 NodeManager 3585 SecondaryNameNode 3768 ResourceManager 4968 Jps 4858 HRegionServer 3451 DataNode 3355 NameNode 4765 HMaster 4671 HQuorumPeer [localhost ~|s jps 3476 HRegionServer] | logging to /home/acadgild/install/hadoop/hadoop-2.6.5/sbin/mr-jobhistory-daemon.sh start historyserver localhost.localdomain.out 12 localdomain.out 13 localdomain.out 13 localdomain.out 13 localdomain.out 13 localdomain.out 14 localdomain.out 15 localdomain
```

Then we have started **hbase shell** and created table 'bulktable' with column families cf1 and cf2.

```
[acadgild@localhost ~]$ hbase shell

2018-08-18 11:08:47,684 WARN [main] util.NativeCodeLoader: Unable to load native-hadoop library for your platform... using builtin-java classes where applicable

SLF43: Found binding in [jar:file:/home/acadgild/install/hbase/hbase-1.2.6/lib/slf4j-log4j12-1.7.5.jar!/org/slf4j/impl/StaticLoggerBinder.class]

SLF43: Found binding in [jar:file:/home/acadgild/install/hadoop/hadoop-2.6.5/share/hadoop/common/lib/slf4j-log4j12-1.7.5.jar!/org/slf4j/impl/StaticLoggerBinder.class]

SLF43: Found binding in [jar:file:/home/acadgild/install/hadoop/hadoop-2.6.5/share/hadoop/common/lib/slf4j-log4j12-1.7.5.jar!/org/slf4j/impl/StaticLoggerBinder.class]

SLF43: See http://www.slf4j.org/codes.html#multiple_bindings for an explanation.

SLF43: Actual binding is of type [org.slf4j.impl.Log4jLoggerFactory]

HBase Shell; enter 'help<RETURN>' for list of supported commands.

Type "exit<RETURN>" to leave the HBase Shell

Version 1.2.6, rUnknown, Mon May 29 02:25:32 CDT 2017

hbase(main):001:0> create 'bulktable','cf1','cf2'

0 row(s) in 3.3150 seconds

⇒ Hbase::Table - bulktable

hbase(main):002:0> ■
```

Then we have created Hbase directory in local system.

```
[acadgild@localhost ~]$ mkdir Hbase
You have new mail in /var/spool/mail/acadgild
[acadgild@localhost ~]$ cd Hbase
[acadgild@localhost Hbase]$ pwd
/home/acadgild/Hbase
```

Then we have created and edited bulk data.tsv file by using **VI editor**.

```
[acadgild@localhost Hbase]$ vi Hbase/bulk_data.tsv

1 Amit 4

2 girija 3

3 jatin 5

4 swati 3
```

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By using below cat command, we could see that data has been inserted successfully in bulk_data.tsv file.

As data should be present in HDFS, we have to copy bulk_data.tsv file from local to HDFS. So we are creating hbase directory in HDFS via command: hadoop fs -mkdir /hbase.

But as this hbase directory is already present, we are getting below error as hbase directory already exists.

```
[acadgild@localhost rbase]s hadoop fs -mkdir /hbase
18/08/18 11:51:44 WARN util.NativeCodeLoader: Unable to load native-hadoop library for your platform... using builtin-java classes where
applicable
mkdir: /hbase': File exists
```

Then we have copied bulk_data.csv file from local to hbase directory in HDFS.

By using below command, you could see content in bulk_data.tsv file:

hadoop fs -cat /hbase/bulk_data.tsv

```
[acadgild@localhost Hbase]s hadoop fs -put bulk_data.tsv /hbase
18/08/18 11:54:40 WARN util.NativeCodeLoader: Unable to load native-hadoop library for your platform... using builtin-java classes where
applicable
[acadgild@localhost Hbase]s hadoop fs -cat /hbase/bulk_data.tsv
18/08/18 11:55:48 WARN util.NativeCodeLoader: Unable to load native-hadoop library for your platform... using builtin-java classes where
applicable
1 Amit 4
2 girija 3
3 jatin 5
4 swati 3
You have new mail in /var/spool/mail/acadgild
```

Now to import data from HDFS to HBase, we have used following command along with arguments<tablename> and <path of data in HDFS>.

hbase org.apache.hadoop.hbase.mapreduce.lmportTsv - Dimporttsv.columns=HBASE_ROW_KEY,cf1:name,cf2:exp bulktable /hbase/bulk_data.tsv

Here, tablename is bulktable & path of data file is /hbase/bulk_data.tsv.

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Then to check whether data got inserted from HDFS into HBase, we have fetched data from bulktable by using HBase command as **scan 'bulktable'**.

We could see that all data (4 rows) which is present in **bulk_data.tsv** file got inserted into bulktable in HBase successfully.

```
hbase(main):002:0> scan 'bulktable'

ROW

COLUMN+CELL

column=cf1:name, timestamp=1534573770704, value=Amit

column=cf2:exp, timestamp=1534573770704, value=girija

column=cf2:exp, timestamp=1534573770704, value=girija

column=cf2:exp, timestamp=1534573770704, value=3

column=cf1:name, timestamp=1534573770704, value=jatin

column=cf2:exp, timestamp=1534573770704, value=5

column=cf2:exp, timestamp=1534573770704, value=swati

column=cf2:exp, timestamp=1534573770704, value=swati

column=cf2:exp, timestamp=1534573770704, value=3

4 row(s) in 0.4230 seconds
```