

Knowledge representati

Inference in First-Order Logic

- Inference in First-Order Logic is used to deduce new facts or sentences from existing sentences.
- Basic terminologies :
- Substitution: Substitution is a fundamental operation performed on terms and formulas. It
 occurs in all inference systems in first-order logic. The substitution is complex in the presence
 of quantifiers in FOL. If we write F[a/x], so it refers to substitute a constant "a" in place of
 variable "x".
- 2. **Equality**: First-Order logic does not only use predicate and terms for making atomic sentences but also uses another way, which is equality in FOL. For this, we can use **equality symbols** which specify that the two terms refer to the same object.

Example: Brother (John) = Smith.

As in the above example, the object referred by the **Brother (John)** is similar to the object referred by **Smith**. The equality symbol can also be used with negation to represent that two terms are not the same objects.

 \neg (x=y) which is equivalent to x \neq y.

Inference rules of FOL

Universal Generalisation

• Universal generalisation is a valid inference rule which states that if premise P(c) is true for any arbitrary element c in the universe of discourse, then we can have a conclusion as \forall x P(x).

- It can be represented as: P(c)• $\forall x P(x)$.
- This rule can be used if we want to show that every element has a similar property.
- In this rule, x must not appear as a free variable.
- **Example:**

Let's represent, P(c): "A byte contains 8 bits", so for \forall x P(x) "All bytes contain 8 bits.", it will also be true.

Universal Elimination

- Universal instantiation is also called as universal elimination or UI is a valid inference rule. It can be applied multiple times to add new sentences.
- · As per UI, we can infer any sentence obtained by substituting a ground term for the variable.
- The UI rule state that we can infer any sentence P(c) by substituting a ground term c (a constant within domain x) from \forall x P(x) for any object in the universe of discourse.

- It can be represented as: P(c)
- Example:

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IF "Every person like ice-cream"=> \forall x P(x) so we can infer that "John likes ice-cream" => P(c)
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Existential Elimination

- It can be applied only once to replace the existential sentence.
- The new KB is not logically equivalent to old KB, but it will be satisfiable if old KB was satisfiable.
- This rule states that one can infer P(c) from the formula given in the form of $\exists x P(x)$ for a new constant symbol c.
- The restriction with this rule is that c used in the rule must be a new term for which P(c) is true.

- It can be represented as: P(c)
- Example:

From the given sentence: $\exists x \text{ Crown}(x) \land \text{OnHead}(x, \text{John})$,

So we can infer: **Crown(K)** \(\text{OnHead(K, John),} \) as long as K does not appear in the knowledge base.

- The above used K is a constant symbol, which is called **Skolem constant**.
- The Existential instantiation is a special case of **Skolemization process**.

Existential Introduction/Generalization

- An existential introduction is also known as an existential generalization, which is a valid inference rule in first-order logic.
- This rule states that if there is some element c in the universe of discourse which has a
 property P, then we can infer that there exists something in the universe which has the
 property P.

- It can be represented as: $\exists x P(x)$
- Example: Let's say that,

"Priyanka got good marks in English."

"Therefore, someone got good marks in English."