

# **Programming with Recursion**

# Recursive Function Call

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- A **recursive call** is a function call in which the called function is the same as the one making the call.
- In other words, *recursion occurs when a function calls itself!*
- We must avoid making an infinite sequence of function calls (infinite recursion).

# Finding a Recursive Solution

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- Each successive recursive call should bring you closer to a situation in which the answer is known.
- A case for which the answer is known (and can be expressed without recursion) is called a **base case**.
- Each recursive algorithm must have at least one base case, as well as the **general** (recursive) case.

# General format for many recursive functions

```
if (some condition for which answer is known)
    // base case
    solution statement
else
    // general case
    recursive function call
```

**SOME EXAMPLES . . .**

# Writing a recursive function to find n factorial

## DISCUSSION

The function call `Factorial(4)` should have value 24, because that is  $4 * 3 * 2 * 1$ .

For a situation in which the answer is known, the value of  $0!$  is 1.

So our **base case** could be along the lines of

```
if ( number == 0 )  
    return 1;
```

# Writing a recursive function to find Factorial(n)

Now for the **general case** . . .

The value of Factorial(n) can be written as  
 $n$  \* the product of the numbers from  $(n - 1)$  to 1,  
that is,

$$n * (n - 1) * (n - 2) * (n - 3) * \dots * 3 * 2 * 1$$


or,  $n! = n * (n - 1)! = n * (n - 1) * (n - 2)! = \dots = \dots 1! = \dots 1 * 0!$

And notice that the recursive call Factorial( $n - 1$ ) gets us  
“closer” to the base case of Factorial(0).

# Recursive Solution

```
public static int Factorial ( int  n )  
// Pre: number is assigned and number >= 0.  
{  
    if ( n == 0) // base case, for termination/stopping  
        return 1 ;  
    else // general case, for recursion  
        return n * Factorial ( n - 1 ) ;  
}
```

# Three-Question Method of verifying recursive functions

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- **Base-Case Question:** Is there a nonrecursive way out of the function?  
The answer should be “yes”
- **Smaller-Caller Question:** Does each recursive function call involve a smaller case of the original problem leading to the base case?  
The answer should be “yes”
- **General-Case Question:** Assuming each recursive call works correctly, does the whole function work correctly?  
The answer should be “yes”



# Another example where recursion comes naturally

- From mathematics, we know that

$$2^0 = 1 \quad \text{and} \quad 2^5 = 2 * 2^4$$

- In general,

$$x^0 = 1 \quad \text{and} \quad x^n = x * x^{n-1} = x * x * x^{n-2}$$

for integer  $x$ , and integer  $n > 0$ .

- Here we are defining  $x^n$  recursively, in terms of  $x^{n-1}$
- $x^n = x * x * x * x * \dots n \text{ times}$

### **// Recursive definition of power function**

```
public static int Power ( int  x,  int  n )
```

```
    // Pre:   n >= 0.  x, n are not both zero
```

```
    // Post:  Function value = x raised to the power n.
```

```
{
    if ( n == 0 )
        return 1;  // base case, for termination/stopping
    else           // general case, for recursion
        return ( x * Power ( x , n-1 ) ) ;
}
```

**Of course, an alternative would have been to use looping instead of a recursive call in the function body.**

# How recursion works?

Power:  $x^n$

$$\begin{aligned}x^4 &= x \cdot x^3 = x \cdot (x \cdot x^2) = x \cdot (x \cdot (x \cdot x^1)) = x \cdot (x \cdot (x \cdot (x \cdot x^0))) \\&= x \cdot (x \cdot (x \cdot (x \cdot 1))) = x \cdot (x \cdot (x \cdot (x))) = x \cdot (x \cdot (x \cdot x)) \\&= x \cdot (x \cdot x \cdot x) = x \cdot x \cdot x \cdot x\end{aligned}$$

call 1	$x^4 = x \cdot x^3$	$= x \cdot x \cdot x \cdot x$
call 2	$x \cdot x^2$	$= x \cdot x \cdot x$
call 3	$x \cdot x^1$	$= x \cdot x$
call 4	$x \cdot x^0$	$= x \cdot 1 = x$
call 5	1	

or alternatively, as

call 1	<code>power(x, 4)</code>	
call 2	<code>power(x, 3)</code>	
call 3	<code>power(x, 2)</code>	
call 4	<code>power(x, 1)</code>	
call 5	<code>power(x, 0)</code>	
call 5	1	
call 4	$x$	
call 3	$x \cdot x$	
call 2	$x \cdot x \cdot x$	
call 1	$x \cdot x \cdot x \cdot x$	

# How recursion works?

Power:  $x^n$

```
static public void main(String args[]) {  
    { ...  
/* 136 */    y = power(5.6,2);  
    ...  
}
```

A trace of the recursive calls is relatively simple, as indicated by this diagram

call 1	power(5.6,2)
call 2	power(5.6,1)
call 3	power(5.6,0)
call 3	1
call 2	5.6
call 1	31.36

# Non-recursive solution

**Power:  $x^n$**

// Pseudo code. Remember:  $x^n = x * x * x * x * \dots n \text{ times} = 1 * x * x * \dots x \text{ n times}$

```
NonRecursivePower(x,n)
{
    Result =1;
    For i = 1 to n
        result = result * x
}
```

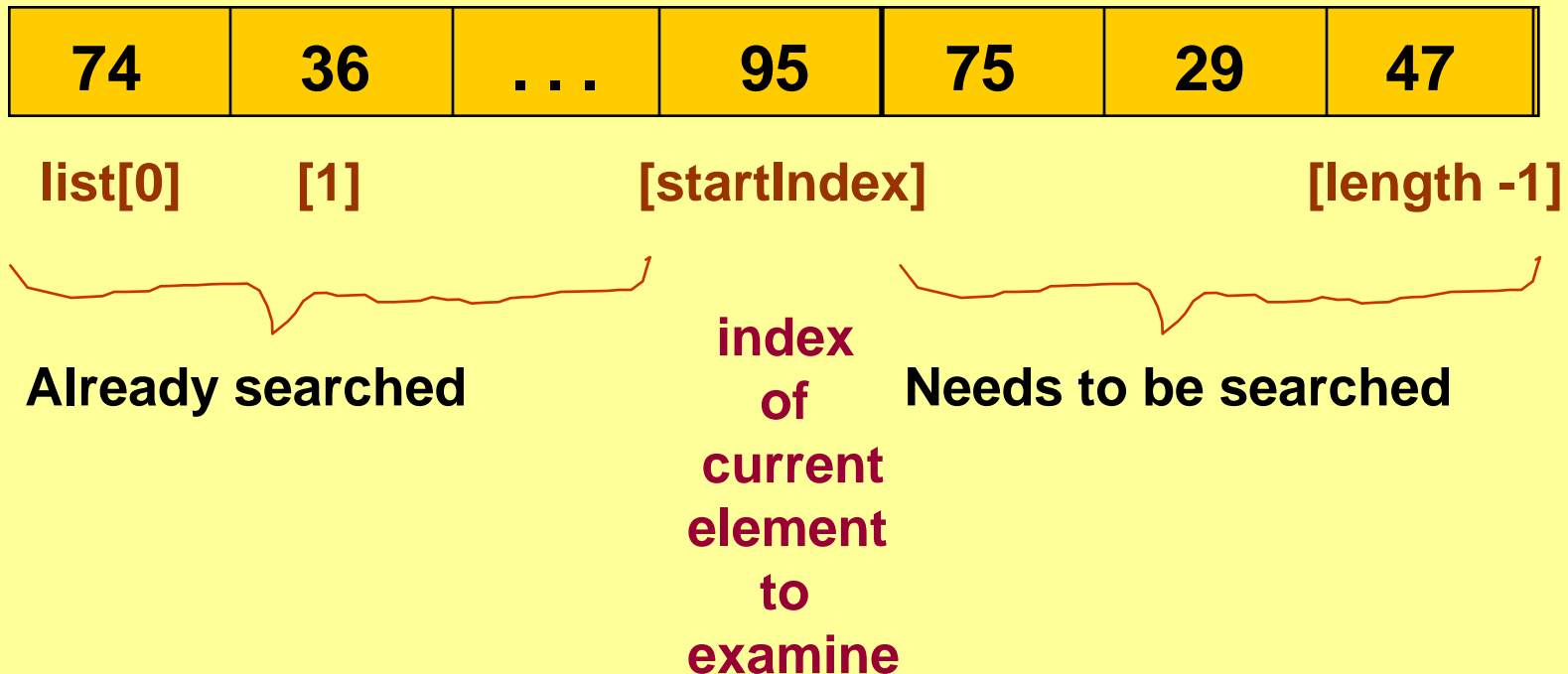
**Jave code**

```
double nonRecPower(double x, int n) {
    double result = 1;
    if (n > 0)
        for (result = x; n > 1; --n)
            result *= x;
    return result;
}
```

# Recursive function to determine if value is in list

## PROTOTYPE

```
public boolean ValueInList ( int list[ ], int value , int startIndex )
```



```

public boolean ValueInList (int list[ ] , int value , int startIndex )

// Searches list for value between positions startIndex
// and list.length-1
// Pre: list [ startIndex ] . . list [ list.length - 1 ]
// contain values to be searched
// Post: Function value =
// ( value exists in list [ startIndex ] . . list [ list.length - 1 ] )
{
    if ( list[startIndex] == value )                // one base case
        return true ;
    else if (startIndex == list.length -1 )          // another base case
        return false ;
    else                                              // general case
        return ValueInList( list, value, startIndex + 1 ) ;
}

main()
{
    ValueInList(list, value, 0);
}

```

# **“Why use recursion?”**

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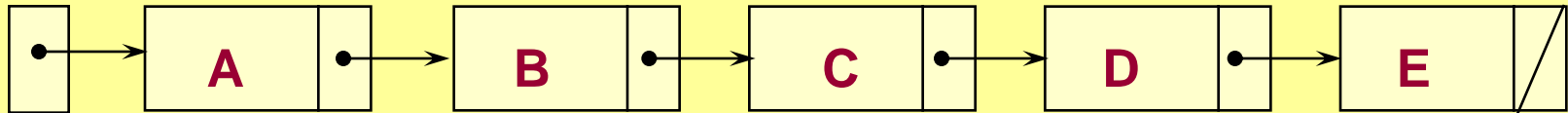
**Those examples could have been written without recursion, using iteration instead. The iterative solution uses a loop, and the recursive solution uses an if statement.**

- However, for certain problems the recursive solution is the most natural solution. This often occurs when pointer variables are used.**
- Recursion is easy to code.**



# Example: Print link list data in reverse order

Head



**THEN, print  
this element**

**FIRST, print out this section of list, backwards**

**So, the output is: E D C B A**

# **Example Continued**

## **Base Case and General Case**

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**A base case may be a solution in terms of a “smaller” list. Certainly for a list with 0 elements, there is no more processing to do.**

**Our general case needs to bring us closer to the base case situation. That is, the number of list elements to be processed decreases by 1 with each recursive call. By printing one element in the general case, and also processing the smaller remaining list, we will eventually reach the situation where 0 list elements are left to be processed.**

**In the general case, we will print the elements of the smaller remaining list in reverse order, and then print the current pointed to element.**

# Example Continued: Solution using recursion

```
public static void reversePrint (LN  pointer)
{
    if (pointer == null)
        return;
    else {
        reversePrint(pointer.next);
        System.out.println(pointer.value);
    };
}

main()
{
    reversePrint(Head);
}
```

# Function `BinarySearch ( )`

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- `BinarySearch` takes **sorted** array `info`, and two subscripts, `fromLoc` and `toLoc`, and `item` as arguments. It returns `false` if `item` is not found in the elements `info[fromLoc...toLoc]`. Otherwise, it returns `true`.
- `BinarySearch` can be written using iteration, or using recursion.

```
found = BinarySearch(info, 25, 0, 14);
```

item fromLoc toLoc

indexes

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
---	---	---	---	---	---	---	---	---	---	----	----	----	----	----

info

0	2	4	6	8	10	12	14	16	18	20	22	24	26	28
---	---	---	---	---	----	----	----	----	----	----	----	----	----	----

16	18	20	22	24	26	28
----	----	----	----	----	----	----

24	26	28
----	----	----

24
----

NOTE:  denotes element examined

## // Recursive definition

```
public boolean BinarySearch (int info[ ] , int item , int fromLoc , int toLoc )
```

```
    // Pre: info [ fromLoc . . toLoc ] sorted in ascending order
```

```
    // Post: Function value = ( item in info [ fromLoc . . toLoc] )
```

```
{    int mid ;
    if ( fromLoc > toLoc )                // base case -- not found
        return false ;
    else {
        mid = ( fromLoc + toLoc ) / 2 ;
        if ( info [ mid ] == item )      // base case-- found at mid
            return true ;
        else if ( item < info [ mid ] )   // search lower half
            return BinarySearch ( info, item, fromLoc, mid-1 ) ;
        else                             // search upper half
            return BinarySearch( info, item, mid + 1, toLoc ) ;
    }
}
```

# Tail Recursion

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- The case in which a function contains **only a single recursive call** and it is the **last statement** to be executed in the function.
- Tail recursion can be replaced by iteration to remove recursion from the solution as in the next example.

# Some Examples

## Tail Recursion

```
void tail (int i) {  
    if (i > 0) {  
        System.out.print (i + "");  
        tail(i-1);  
    }  
}
```

## Non-Tail Recursion

```
void nonTail (int i) {  
    if (i > 0) {  
        nonTail(i-1);  
        System.out.print (i + "");  
        nonTail(i-1);  
    }  
}
```

## Iteration/Loop

```
void iterativeEquivalentOfTail (int i) {  
    for ( ; i > 0; i--)  
        System.out.print(i+ "");  
}
```



**Another Example: press keyboard characters, finish by pressing “Enter”, print the characters in reverse order**

**By Non-Tail Recursion**

```
void reverse() {  
    char ch = getChar();  
    if (ch != '\n') {  
        reverse();  
        System.out.print(ch);  
    }  
}
```

# Another Example: Search Value

## By Tail Recursion

```
public boolean ValueInList (int list[ ], int value , int startIndex )  
  
// Searches list for value between positions startIndex  
// and list.length-1  
// Pre: list [ startIndex ] . . list [ list.length - 1 ]  
// contain values to be searched  
// Post: Function value =  
// ( value exists in list [ startIndex ] . . list [ list.length - 1 ] )  
{  
    if ( list[startIndex] == value )                // one base case  
        return true ;  
    else if (startIndex == list.length -1 )          // another base case  
        return false ;  
    else                                             // general case  
        return ValueInList( list, value, startIndex + 1 ) ;  
}
```

## By Loop/Iteration

### ***// ITERATIVE SOLUTION***

```
public boolean ValueInList (int list[ ] , int value , int startIndex )  
  
// Searches list for value between positions startIndex  
// and list.length-1  
// Pre: list.info[ startIndex ] . . list.info[ list.length - 1 ]  
// contain values to be searched  
// Post: Function value =  
// ( value exists in list.info[ startIndex ] . . list.info[ list.length - 1 ] )  
{ boolean found;  
  found = false ;  
  while ( !found && startIndex < list.length )  
  {    if ( value == list [ startIndex ] )  
        found = true ;  
      else    startIndex++ ;  
  }  
  return found ;  
}
```

# Recursion is not always good

- It may be very slow when excessive recursion calls
- Example: Compute Fibonacci Number

$$\text{Fib}(n) = \begin{cases} n & \text{if } n < 2 \\ \text{Fib}(n-2) + \text{Fib}(n-1) & \text{otherwise} \end{cases}$$

0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, . . .

# Compute Fibonacci Number: Continue...

## Recursive solution:

```
int Fib (int n) {  
    if (n < 2)  
        return n;  
    else return Fib(n-2) + Fib(n-1);  
}
```

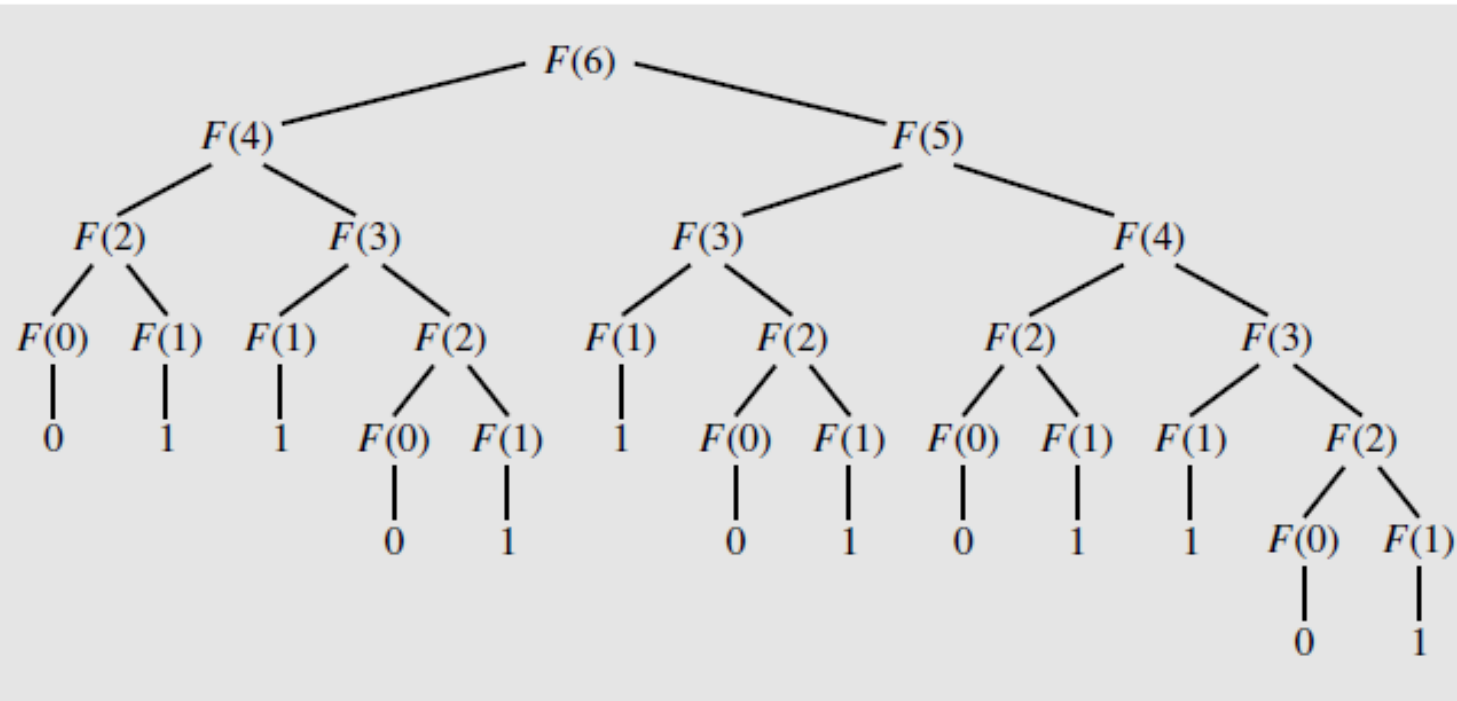
Fib(6) =	Fib(4)	+ Fib(5)
= Fib(2)	+ Fib(3)	+ Fib(5)
= Fib(0)+Fib(1)	+ Fib(3)	+ Fib(5)
= 0 + 1	+ Fib(3)	+ Fib(5)
= 1	+ Fib(1)+ Fib(2)	+ Fib(5)
= 1	+ Fib(1)+Fib(0)+Fib(1)	+ Fib(5)

etc.

# Compute Fibonacci Number: Continue...

## Recursive solution:

The tree of calls for `Fib(6)`.



- Do you see what is the problem here?
- 25 calls !!! to compute  $F(6)$
- The problem is: Same function is repeated again and again.  
For example,  $F(1)$  has been computed 8 times!

# Compute Fibonacci Number: Continue...

## Recursive solution:

Number of addition operations and number of recursive calls to calculate Fibonacci numbers.

n	Fib(n+1)	Number of Additions	Number of Calls
6	13	12	25
10	89	88	177
15	987	986	1,973
20	10,946	10,945	21,891
25	121,393	121,392	242,785
30	1,346,269	1,346,268	2,692,537

# Compute Fibonacci Number: Continue...

## Iterative solution:

An iterative algorithm may be produced rather easily as follows:

```
int iterativeFib (int n) {  
    if (n < 2)  
        return n;  
    else {  
        int i = 2, tmp, current = 1, last = 0;  
  
        for ( ; i <= n; ++i) {  
            tmp = current;  
            current += last;  
            last = tmp;  
        }  
        return current;  
    }  
}
```



# Compute Fibonacci Number: Continue...

## Iterative Solution Vs Recursion

Comparison of iterative and recursive algorithms for calculating Fibonacci numbers.

n	Number of Additions	Assignments	
		Iterative Algorithm	Recursive Algorithm
6	5	15	25
10	9	27	177
15	14	42	1,973
20	19	57	21,891
25	24	72	242,785
30	29	87	2,692,537

**Non-recursive solution**



# Use a recursive solution when:

- The depth of recursive calls is relatively “shallow” compared to the size of the problem.
- The recursive version does about the same amount of work as the nonrecursive version.
- The recursive version is shorter and simpler than the nonrecursive solution.

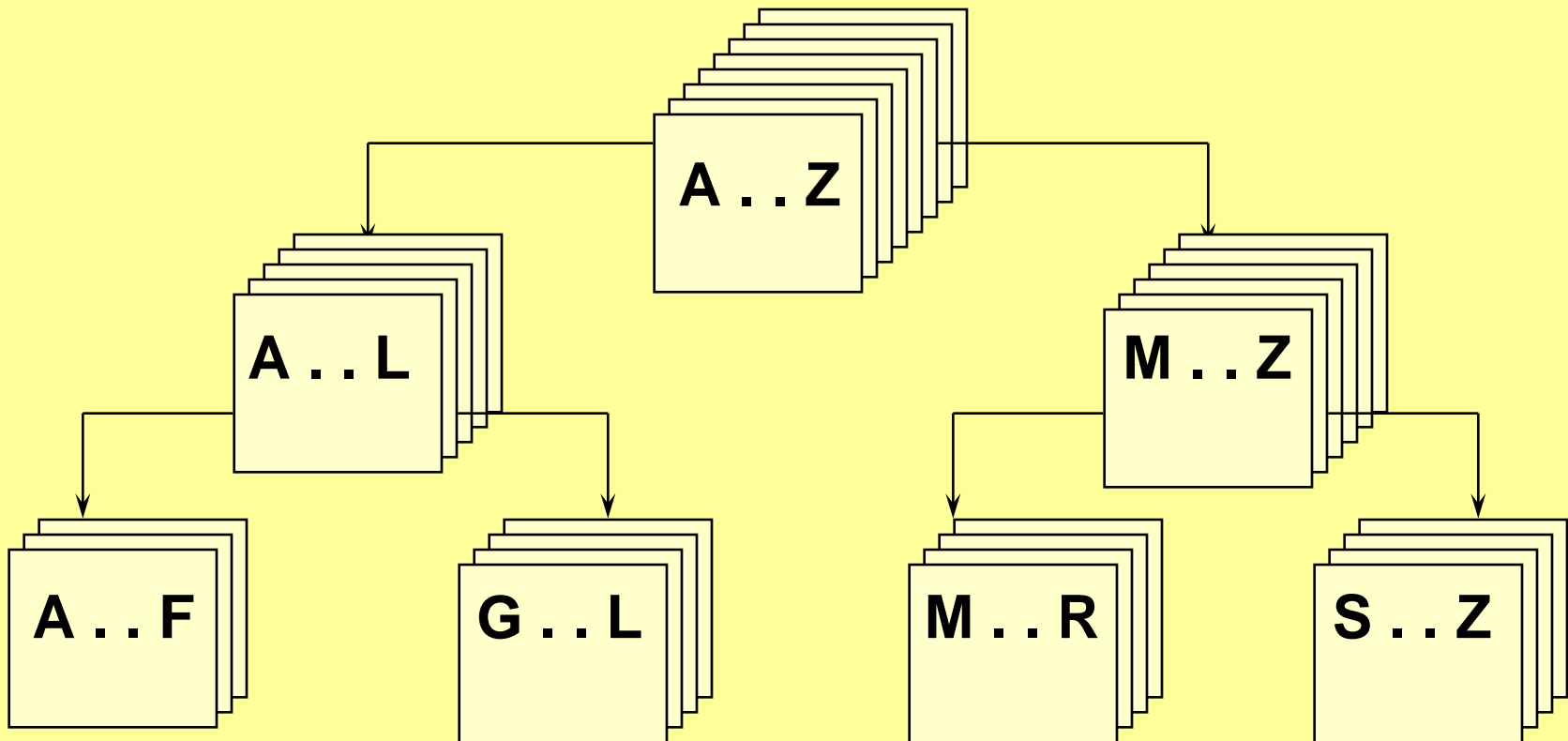
**SHALLOW DEPTH**

**EFFICIENCY**

**CLARITY**

# Example: Quick sort

- Simple and recursive, but fastest sorting algorithm



# Before call to function Split

**splitVal = 9**

**GOAL: place splitVal in its proper position with  
all values less than or equal to splitVal on its left  
and all larger values on its right**

<b>9</b>	<b>20</b>	<b>6</b>	<b>10</b>	<b>14</b>	<b>3</b>	<b>60</b>	<b>11</b>
----------	-----------	----------	-----------	-----------	----------	-----------	-----------

**values[first]**

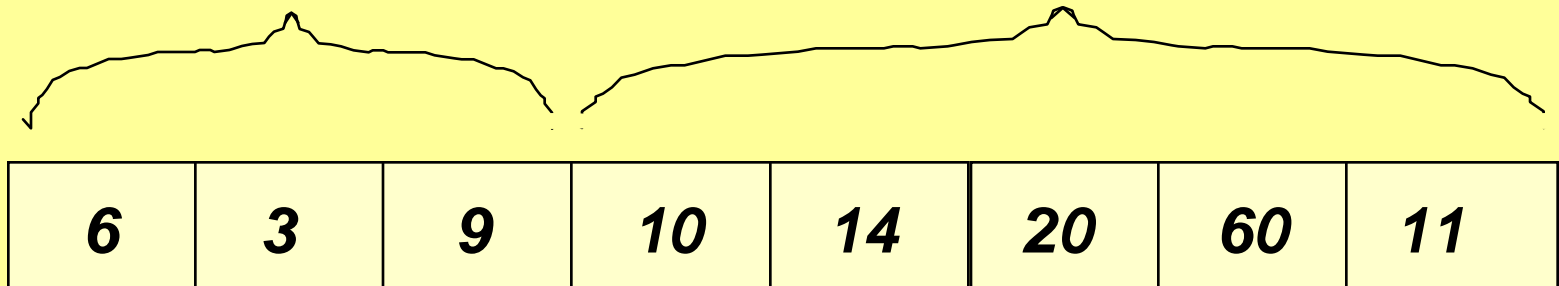
**[last]**

# After call to function Split

**splitVal = 9**

**smaller values**

**larger values**



**values[first]**

**[splitPoint]**

**[last]**

**// Recursive quick sort algorithm**

**public void QuickSort (int values[ ], int first , int last )**

**// Pre: first <= last**

**// Post: Sorts array values[ first . . last ] into ascending order**

**{**

**if ( first < last )** ***// general case***

**{ int splitPoint ;**

**Split ( values, first, last, splitPoint ) ;**

***// values [ first ] . . values[splitPoint - 1 ] <= splitVal***

***// values [ splitPoint ] = splitVal***

***// values [ splitPoint + 1 ] . . values[ last ] > splitVal***

**QuickSort( values, first, splitPoint - 1 ) ;**

**QuickSort( values, splitPoint + 1, last );**

**}**

**};**

***Remember: Quick Sort is the best sorting algorithm***