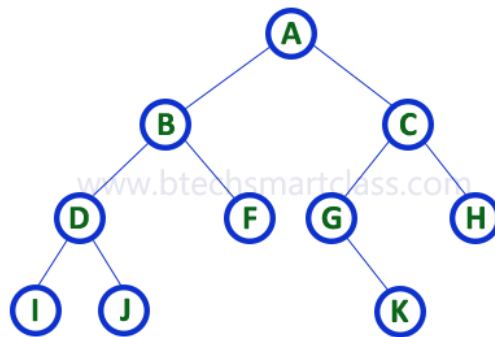


## Binary Tree Representations

A binary tree data structure is represented using two methods. Those methods are as follows...

- **Array Representation**
- **Linked List Representation**

Consider the following binary tree...



### 1. Array Representation

In array representation of binary tree, we use a one dimensional array (1-D Array) to represent a binary tree.

Consider the above example of binary tree and it is represented as follows...



To represent a binary tree of depth ' $n$ ' using array representation, we need one dimensional array with a maximum size of  $2^{n+1} - 1$ .

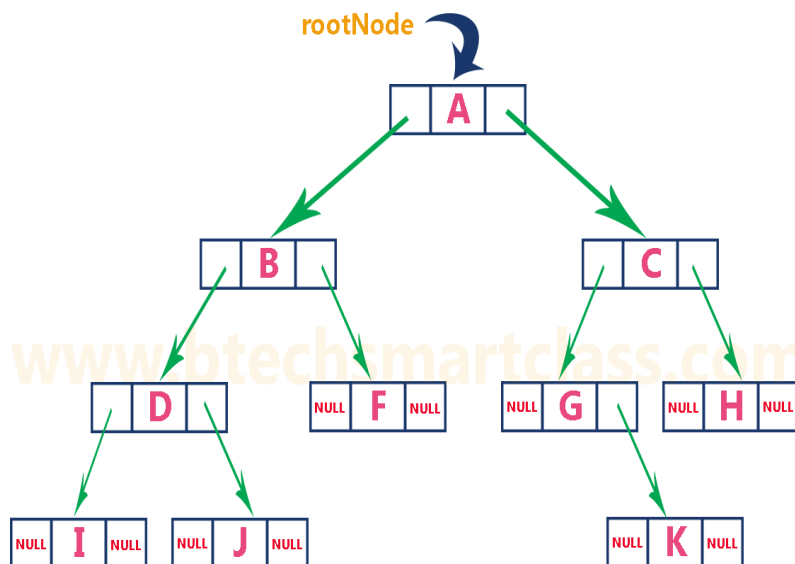
### 2. Linked List Representation

We use double linked list to represent a binary tree. In a double linked list, every node consists of three fields. First field for storing left child address, second for storing actual data and third for storing right child address.

In this linked list representation, a node has the following structure...

<b>Left Child</b> Address	<b>Data</b>	<b>Right Child</b> Address
------------------------------	-------------	-------------------------------

The above example of binary tree represented using Linked list representation is shown as follows...



### Tree traversal techniques:

Traversal is a process to visit all the nodes of a tree and may print their values too. Because, all nodes are connected via edges (links) we always start from the root (head) node. That is, we cannot randomly access a node in a tree. There are three ways which we use to traverse a tree –

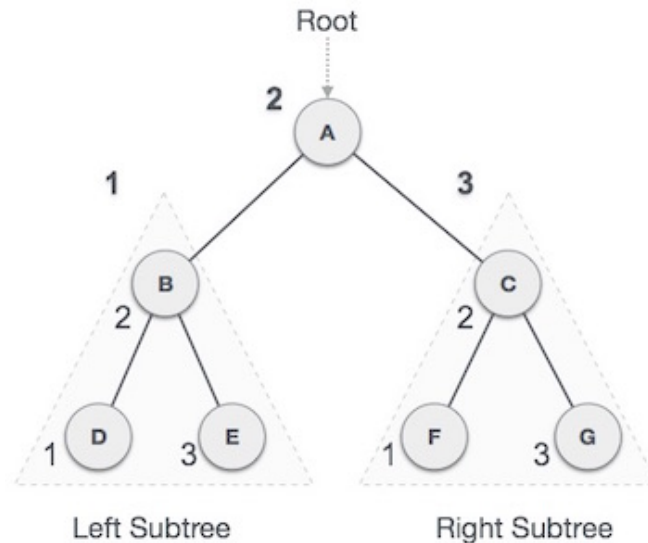
- In-order Traversal
- Pre-order Traversal
- Post-order Traversal

Generally, we traverse a tree to search or locate a given item or key in the tree or to print all the values it contains.

### **In-order Traversal**

In this traversal method, the left subtree is visited first, then the root and later the right sub-tree. We should always remember that every node may represent a subtree itself.

If a binary tree is traversed **in-order**, the output will produce sorted key values in an ascending order.



We start from **A**, and following in-order traversal, we move to its left subtree **B**. **B** is also traversed in-order. The process goes on until all the nodes are visited. The output of inorder traversal of this tree will be –

$D \rightarrow B \rightarrow E \rightarrow A \rightarrow F \rightarrow C \rightarrow G$

### Algorithm

Until all nodes are traversed –

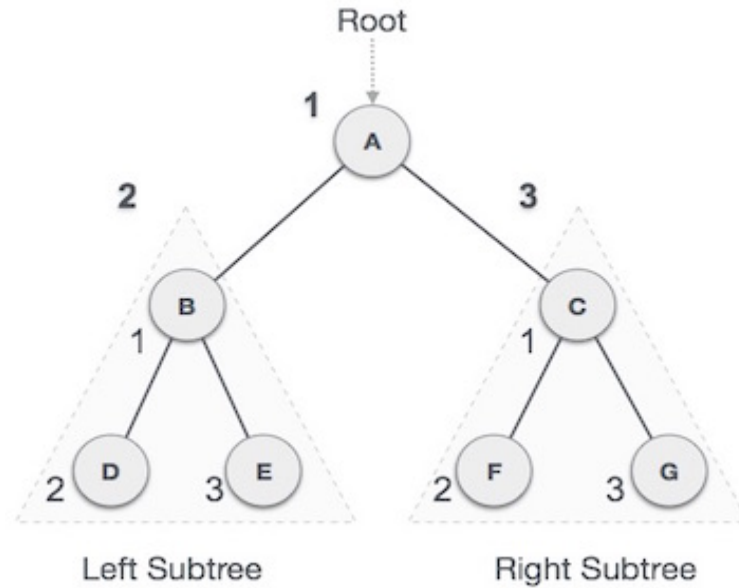
**Step 1** – Recursively traverse left subtree.

**Step 2** – Visit root node.

**Step 3** – Recursively traverse right subtree.

### Pre-order Traversal

In this traversal method, the root node is visited first, then the left subtree and finally the right subtree.



We start from **A**, and following pre-order traversal, we first visit **A** itself and then move to its left subtree **B**. **B** is also traversed pre-order. The process goes on until all the nodes are visited. The output of pre-order traversal of this tree will be –

**$A \rightarrow B \rightarrow D \rightarrow E \rightarrow C \rightarrow F \rightarrow G$**

### Algorithm

Until all nodes are traversed –

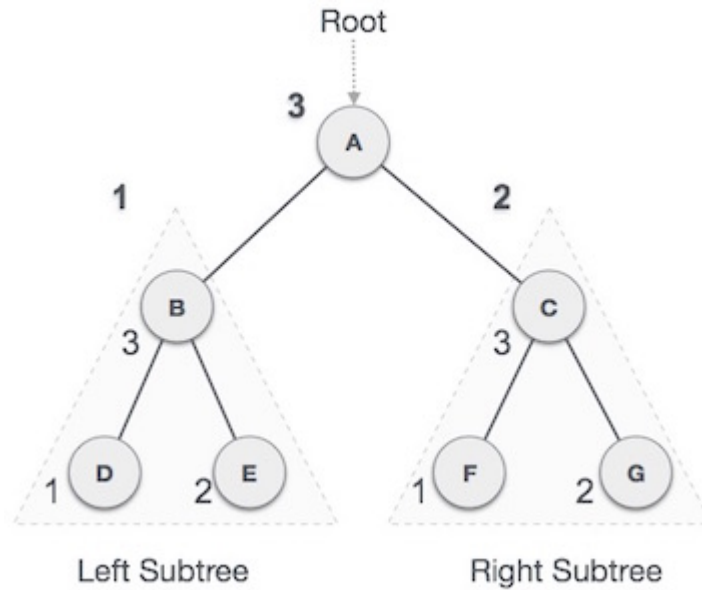
**Step 1** – Visit root node.

**Step 2** – Recursively traverse left subtree.

**Step 3** – Recursively traverse right subtree.

### Post-order Traversal

In this traversal method, the root node is visited last, hence the name. First we traverse the left subtree, then the right subtree and finally the root node.



We start from **A**, and following Post-order traversal, we first visit the left subtree **B**. **B** is also traversed post-order. The process goes on until all the nodes are visited. The output of post-order traversal of this tree will be –

**$D \rightarrow E \rightarrow B \rightarrow F \rightarrow G \rightarrow C \rightarrow A$**

### Algorithm

Until all nodes are traversed –

**Step 1** – Recursively traverse left subtree.

**Step 2** – Recursively traverse right subtree.

**Step 3** – Visit root node.