CSED 2014

Hashing
Equivalent Classes

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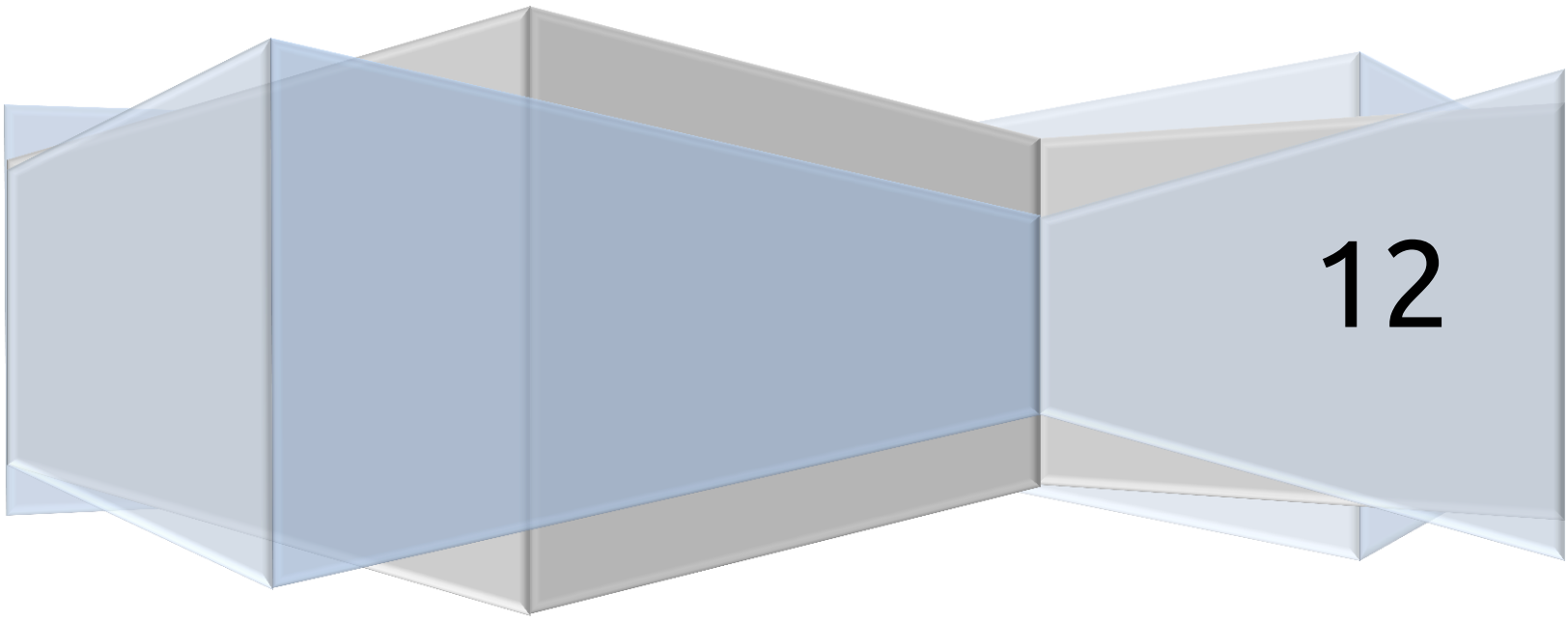
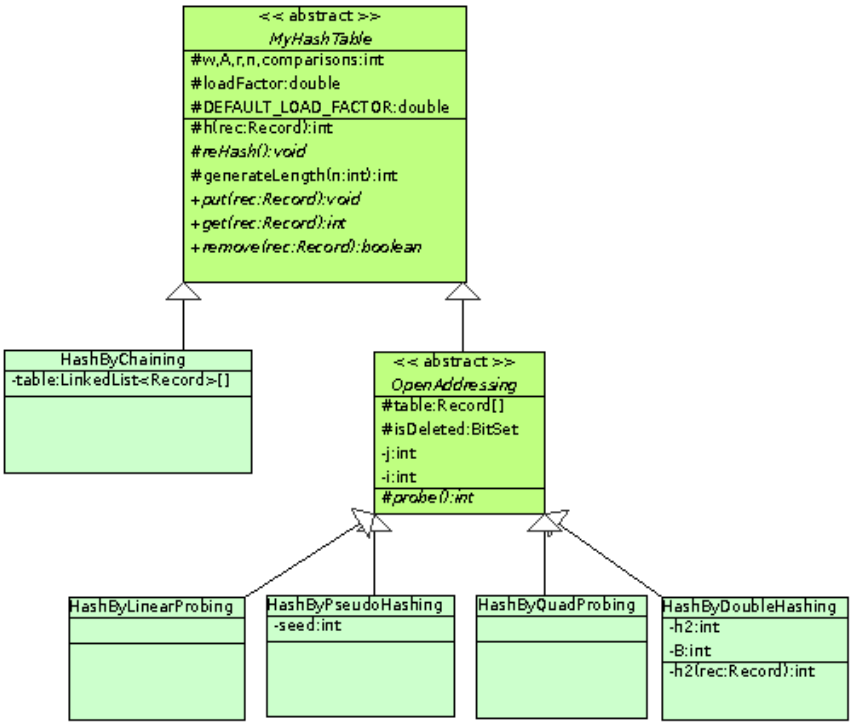


UML :



Algotithms :

Hashing algorithms:

h1(Record rec) //primary hashing function based on name field
k <- rec.getName().hashCode();
//A*k mod 2^w >> (w-r)
//taking the mod to 2^w = taking the w bits from the number
return A*k >> (w-r);

h2(Record rec) //secondary hashing used in DoubleHashing
k <- rec.getName().hashCode();
h2 = (A*k + B)% (1<<r);
//where A and B are Primes
return h2;

Algorithm : ReHash()

```
make a new table with double size
re-put all entries in the old table
```

Hash By Chaning

```
Algorithm : put (Record rec)
        index <- h(rec);
        add record to table at index ;
        comparisons <- 0;
        loadFactor <- ++n / (1<<r) ;
        if (loadFactor > DEFAULT_LOAD_FACTOR)
                reHash() ;

Algorithm : get(Record rec)
        comparisons <- 0;
        index <- h(rec);
                for (all entries at index in table){
                        comparisons++;
                        if (rec == entry)return index;
                }
        //if didn't return at the for loop it means that rec is not found in the table;
        return comparisons =-1 ;

Algorithm : Remove(Record rec)
        index = h(rec);
        comparisons = 0;

                for (all entries at index in table){
                        if (rec== entry)
                                comparisons++;
                                remove rec from table;
                                return true;
                }
                //if didn't return in the for loop then rec was not found in the table
                comparisons = -1;
                return false ;
```

Open Addressing Class

```
Algorithm : put(Record rec)
        h <- h(rec);
        index <- h;
        j = comparisons <- 0 ; // j is the number of collisions
```

```
        while(table[index]!=null  &&  record at index is not marked for deletion){
                index <- probe(h,j++);
                comparisons++;
        }
        table[index] <- rec ;

        loadFactor <- ++n / (1<<r) ;
        if (loadFactor > DEFAULT_LOAD_FACTOR)
                reHash() ;

Algorithms : get(Record rec)
        // To improve, when an element is searched and found in the table,
        // the element is relocated to the first location marked for deletion
        // that was probed during the search.

        h = h(rec);
        int i <- h;
        j == comparisons <- 0 ;
        int firstDeleted <- -1;
        boolean found = false ;

        //don't have to check that # of probes < length  as i rehash when the load
factor > 0.7
        // for sure i'll find an empty slot

                while(table[i] != null){
                                comparisons++;
                                if (record was marked for deletion)
                                        if (firstDeleted == -1)
                                                firstDeleted = i;
                                        i <- probe(h,j++);

                                else
                                        found = rec.equals(table[i]);
                                        // update the table
                                        if (found && firstDeleted != -1)
                                                table[firstDeleted] = rec ;
                                        break;

                }

                if (found)
                        return i;
                else
```

return comparisons = -1;

Algorithm : remove(Record rec)
index <- get(rec);
if (found)
return false;

mark entry as deleted
n--;
return true;

Linear Probing Class

Algorithm : probe(int i, int j)
return (i+j)%(1<<r);

Quad Probing Class

Algorithm : probe(int i, int j)
return (i+j*j)%(1<<r);

Double Hashig Class

Algorithm : probe(int i, int j)
return h+j*h2;

pseudo Random Probing Class

Algorithm : probe(int i, int j)
return (seed*(h+j) + j*7549)%(1<<r);

DataStrucures :

BitSet : to mark entries in the table as deleted (in open addressing classes)

Order comparison

hashing by external Chaining
memory : n records + n links (int size links)
order : \theta (1 + \alpha) == O(m)
where m is the table size

hashing by open Addressing
memory : n records
order : \theta (1/(1-\alpha))