MP4: Page Manager II

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Assigned Tasks

• Part 1: Support for Large Address Spaces - Completed.

I have utilized the recursive page table lookup method to map the logical address space to the physical frames. When a logical address is in the following format:

X:10Y:10 offset: 12

The page directory can be accessed as follows:

1023:10 1023:10

page directory: 10

offset : 2

Similarly, the page table can be accessed as follows:

1023:10

page directory: 10

page table: 10

offset: 2

• Part 2: Preparing class PageTable to handle Virtual Memory Pools - Completed.

The function PageTable::register_pool() establishes a connection between the Page Table and the Virtual Memory Pool by maintaining a linked list of Virtual Memory Pools. This linked list contains information about the regions available for virtual memory allocation, including the base address, size of each region, Page Table, and Frame Pool to which they belong. When a Page Fault occurs in a Page Table, it checks if the fault address falls within the range of the base address and base address + size to ensure its validity in the Virtual Memory Pool to which the Page Table belongs. If the fault address satisfies this condition, then page fault is handled. The Virtual Memory Pool keeps track of regions used for memory allocation, and the Page Table utilizes this information to validate and manage Page Faults.

• Part 3: An Allocator and De-allocator for Virtual Memory - Completed.

Virtual memory allocates and de-allocates memory in integral multiples of pages using VMPool::allocate and VMPool::release respectively in the specific virtual memory pool. I have used the array regions of data type region_data which stores the base address in base_addr and size in size for each region.

System Design

The main objective of machine problem 4 is the implementation of a virtual memory manager and a basic virtual memory allocator, with the goal of supporting a large address space. We use a recursive page table lookup approach in this machine problem.

Code Description

During the implementation of this machine problem, I made changes to the following six files:

- 1. cont_frame_pool.H
- $2. \ cont_frame_pool.C$
- 3. page_table.H
- 4. page_table.C
- $5. \text{ vm_pool.H}$
- 6. vm_pool.C

I used the same code for cont_frame_pool.C, cont_frame_pool.H that I implemented for MP2. I made changes to my existing page_table.C and page_table.H that I have implemented in MP3. Additionally, I wrote definitions for the functions defined in vm_pool.H in vm_pool.C.

1. page_table.H

In page_table.H, I have added two static variables vm_pool_head and vm_pool_current both of type VMPool* to store the head and current element of linked list of virtual memory pools. It also includes two functions: PDE_address(addr) returns the address of the Page Directory Entry (PDE) location for a given address, and PTE_address(addr) returns the address of the Page Table Entry (PTE) location for a given address.

Figure 1: page_table.H

2. page_table.C

(a) init_paging: This function is used to initialize static private data members in a class, including the kernel frame pools, memory frame pools, and the shared size for the page table.

Figure 2: **init_paging**

(b) PageTable (Class Constructor): The constructor is used to construct the page table object. It initializes the first page directory by assigning the free frame from the kernel frame pool and marking it as valid (present). It also initializes the first page table by assigning the free frame from the process frame pool. The first directory entry holds the first page table, while all the remaining directories are marked as invalid (not present), except for the last page directory entry, which points to the first page directory entry for the recursive page table lookup. Finally, it initializes variables for virtual memory pool management, such as vm_pool_head and vm_pool_current.

```
PageTable::PageTable()

{
// Constructor for PageTable class
Console::puts("Constructed Page Table object Start\n");
unsigned long page directory frame number * kernel nem pool.-get_frames(1);
unsigned long page directory frame number * kernel nem pool.-get_frames(1);
page_directory = (unsigned long *)[page_directory frame_number * PAGE_SIZE);
// Last pode as valid and pointing to *First pode for recurive table look up
page_directory(IBMIRES PARAGE * 1] * (((unsigned long) page_directory) | $0.3);
unsigned long page_table frame number = process nem_pool.-spet_frames(1);
unsigned long * page_table = (unsigned long *)[page_table_frame_number * PAGE_SIZE);
// mapping the first 4MB of memory
for (unsigned long i = 0, physical_address = 0; i < ENTRIES_PER_PAGE; i++, physical_address +* PAGE_SIZE)

// read/vrite, present(011 in binary)
page_table[i] = physical_address | $0.3);
// attribute set to: supervisor level.
// read/vrite, present(011 in binary)
// first pde as valid pointing to page_table | $0.3);
// page_cirectory(IBMIRES_PER_PAGE - 1; i++)

{
// attribute set to: supervisor level.
// read/vrite, present(011 in binary)
page_directory(0] = (((unsigned long)page_table) | $0.3);
// attribute set to: supervisor level.
// read/vrite, or present(014 in binary)
page_directory(1) = 0 | $0.2;
}

vm pool_head * NULL;
console::puts("Constructed Page Table object End\n");
}
```

Figure 3: PageTable (Class Constructor)

(c) **load**: This function sets the current_page_table pointer to the current instance of the PageTable object. Then, it loads the current page directory into register CR3 using write_cr3(). The page table is then loaded.

```
void PageTable::load()
{
    // Load the page table into the CR3 register
    Console::puts("Loaded page table Start\n");
    current_page_table = this;
    write_cr3((unsigned long)current_page_table->page_directory);
    Console::puts("Loaded page table End\n");
}
```

Figure 4: load

(d) **enable_paging**: This function is used to enable paging by setting the paging bit (bit 31) of CR0 to 1 using read_cr0 to read the contents of CR0 and write_cr0 to write the contents

to CR0. It also sets the boolean paging_enabled to true. Before enabling paging, the page directory and page table should be set up and loaded correctly.

```
void PageTable::enable_paging()
{
    Console::puts("Enabled paging Start\n");
    unsigned long cr0_reg = (unsigned long)(read_cr0() | 0x80000000);
    paging_enabled = 1;
    write_cr0(cr0_reg);
    Console::puts("Enabled paging End\n");
}
```

Figure 5: enable_paging

(e) handle_fault: This function is utilized to handle raised faults. The function retrieves the error code from the provided REGS structure. It examines the error code received from the register. If the least significant bit of the error code is 0, indicating a page fault, the function proceeds to handle the fault. Subsequently, it retrieves the faulty address from the CR2 register using the read_cr02 function. Additionally, it reads the current page directory list and directory location from CR3 using the read_cr3 function and the faulty address, respectively. It checks whether the faulty address is a legitimate virtual memory address by iterating over the virtual memory pool list. If it is not legitimate, then the execution is halted with a message, 'Not a legitimate address'. If the page directory entry for the directory location is not found, indicating a problem with the directory, the function creates a page table and allocates a new page table from the process memory pool. It updates the page directory entry accordingly and sets all the page table entries as not being present. If the page directory entry is found, indicating an existing issue with the page table, the function allocates a frame from the process memory pool and updates the corresponding page table entry. In the event of a different scenario than the aforementioned, the function halts execution with a message stating, 'Something went wrong'.

```
void PageTable::handle_fault(REGS *_r)
{
    // Handle page faults
    Console::puts('handle fault Start\n');
    console::puts('handle fault exact recovered;
    if (fer_code & 601) = 0.00
}    console::puts('handle fault err occuured\n');
    unsigned long faulty_address = (unsigned long);
    unsigned long faulty_address = (unsigned long *)(read_cr3());
    unsigned long directory_location = (faulty_address) >> 22;
    brook_interator;
    for (iterator = vm_pool_head; iterator != NULL; iterator = iterator-next)
    {
        bool temp_bool = iterator-is_legitimate(faulty_address) == true;
        if (temp_bool)
        if (temp_bool)
        is_legitimate_vm_address = true;
        break;
    }
}

if ((lpage_directory_list(directory_location) & 0x1) == 0x0)
{
        Console::puts('directory issue and new page table*);
        Console::puts('directory issue and new page table of the page_table of the page_ta
```

Figure 6: handle_fault

(f) **register_pool**: This function is used to register a VMPool instance into a linked list of virtual memory pools. It assigns the instance to the head of the linked list of virtual memory pools if the list is empty or appends it to the end of the list if it's not.

```
void PageTable::register_pool(WPool *_vm_pool)
{
    // Register a VMPool instance
    Console::puts("registered VM pool - start\n");
    if \univm_pool_head = \univm_NLL\)
{
        Console::puts("Empty Head.\n");
        // if list is empty assign it to head
        vm_pool_head = _vm_pool;
    }
    else
    {
        Console::puts("Non Empty Head.\n");
        // if list is not empty make next point to it
        vm_pool_current-next = \univm_pool;
    }
    // in both cases current next should be null and current should be point to cureent _vm_pool
        vm_pool_current-next = \univm_pool;
        vm_pool_current-next = \univm_pool
        vm_pool_cu
```

Figure 7: register_pool

(g) **free_page**: This function is responsible for freeing a page and flushing the Translation Lookaside Buffer (TLB). It determines the frame number associated with the page based on its address and invokes the corresponding **release_frame** function from the process pool to release that frame. Additionally, it marks the page as not present in the page table entry. Finally, it ensures that the TLB is updated to reflect the changes in memory mapping.

```
void PageTable::free page(unipned long_page_no)
{
   // Free a page and flush the TLB
   Console::psts(*freed page - start\n');
   unisped long page_directory_location * PDE_address(_page_no);
   unisped long page_directory_clocation * PDE_address(_page_no);
   unisped long frace_no = (page_directory_clocation);
   page_directory_clocation * PAGE_SIZE);
   process_nee_pool.orclease_frames(frame_no);
   page_directory_clocation * page_directory_clocation * page_directory_clocation * PAGE_SIZE);
   page_directory_clocation * PAGE_SIZE);
   long();
   Console::puts(*freed page - end\n');
}
```

Figure 8: **free_page**

(h) PDE_address: This function returns the Page Directory Entry (PDE) Location.

```
// Return the address of the Page Directory Entry (PDE) Location
unsigned long PageTable::PDE_address(unsigned long addr)
{
   unsigned long page_directory_location = (addr & 0xFFC000000) >> 22;
   return page_directory_location;
}
```

Figure 9: PDE_address

(i) PTE_address: This function returns the Page Table Entry (PTE) Location.

```
// Return the address of the Page Table Entry (PTE) Location
vunsigned long PageTable::PTE_address(unsigned long addr)
{
   unsigned long page_table_location = (addr & 0x003FF000) >> 12;
   return page_table_location;
}
```

Figure 10: PTE_address

3. vm_pool.H

In vm_pool.H, a data type called region_data is defined with attributes base_addr and size for specifying the base address and size of each region in a virtual memory pool. Other variables include base_address, size, frame_pool, page_table, and a list of memory pools using region_data. There are variables for storing the available virtual memory size in available_size and the total number of regions in total_count. The next attribute of data type VMPool * is used to store the address of the next virtual memory pool region and is initialized as NULL.

```
class VMPool

/* Virtual Memory Pool */
You, 3 days ago * method definitions for vm pool added
private:
/* -- DEFINE YOUR VIRTUAL MEMORY POOL DATA STRUCTURE(s) HERE. */
You, yesterday|1 author (You)
class region_data
{
  public:
    unsigned long base_addr;
    unsigned long size;
};

unsigned long size;
contFramePool *frame_pool;
PageTable *page_table;
region_data *regions;
unsigned long available_size;
unsigned long available_size;
unsigned long total_count;

public:
WMPool *next = NULL;
WMPool (unsigned long base address.)
```

Figure 11: vm_pool.H

4. vm_pool.C

(a) VMPool (Class Constructor): The constructor is used to initialize private data members such as base_address, size, page_table, frame_pool of virtual memory pool in a class. Additionally, I have initialized the pointer to the next virtual memory pool as NULL, available_size to the size of virtual memory pool, total_count to number of region in a virtual memory pool, and reserved the first region for storing the region data.

```
VMPool::VMPool(unsigned long _base_address,
               unsigned long _size,
ContFramePool *_frame_pool,
               PageTable * page table)
   Console::puts("Constructed VMPool object - start.\n");
   base address = base address;
    frame_pool = _frame_pool;
   page_table = _page_table;
   next = NULL;
    available_size = _size;
    total count = \theta;
    page table->register pool(this);
    region data *temp region = (region data *)base address;
    temp_region[0].base_addr = base_address;
    temp region[0].size = PageTable::PAGE SIZE;
   regions = temp_region;
   available size -= PageTable::PAGE SIZE;
    total count += 1:
   Console::puts("Constructed VMPool object - end.\n");
```

Figure 12: VMPool (Class Constructor)

(b) **allocate**: This function checks for sufficient memory and allocates an area of the requested size if available; otherwise, it stops the execution with the message 'No free size available'. The available memory size is updated to reflect the allocation of the remaining size and increments the number of regions. Finally, it returns the location address of the given region.

Figure 13: allocate

(c) release: This function releases a memory area based on the given start address. It first checks whether the given start address is legitimate or not. If it is, the function searches for the start address among all the available used regions. If it is not found, the program stops execution with the message 'No such region found starting with this start address.' If it is found, the function frees the corresponding pages in the page table. The function adjusts the array by shifting subsequent regions, updates the count and available memory size, and finally flushes the TLB. If the address is not legitimate, the function halts execution with the message 'Not Legitimate - start address'.

```
(region relase index < 0)
 unsigned long curr_addr = _start_address;
unsigned long number_of_pages = regions[region_relase_index].size / PageTable::PAGE_SIZE;
for (unsigned long i = 0; i < number_of_pages; i++)
        page_table->free_page(curr_addr);
curr_addr += PageTable::PAGE SIZE;
  unsigned long curr_addr = _start_address;
unsigned long number_of pages = regions[region_relase_index].size / PageTable::PAGE_SIZE;
for (unsigned long i = 0; i < number_of_pages; i++)
        page_table->free_page(curr_addr);
curr_addr += PageTable::PAGE_SIZE;
 total_count -= 1;
available size += regions[total count].size;
 regions[region_relase_index].base_addr = regions[total_count].base_addr;
regions[region_relase_index].size = regions[total_count].size;
```

Figure 14: release

(d) **is_legitimate**: This function checks if a given address is within any allocated memory region. It iterates through the memory regions and compares the address with the base address and size of each region. If a match is found, it returns true, indicating the address is legitimate. If no match is found, it returns false. The function outputs messages indicating the start and end of the address legitimacy check.

Figure 15: is_legitimate

Testing

During the development of the code, I wrote several Console:puts() and Console:putui() statements to pinpoint where my code was breaking and to understand if the logic was incorrect or not performing as expected. During testing, the program was running infinitely. Additionally, test case 0 was failing. After making necessary changes and analyzing the console debug logs, the code ran successfully. I removed all Console statements and reverted Kernel.C to its original state. Initially, the page table implementation was tested to ensure its functionality. Subsequently, the macro known as _TEST_PAGE_TABLE_ was commented for testing the virtual memory (VM) pools. The following screenshots illustrate this process.

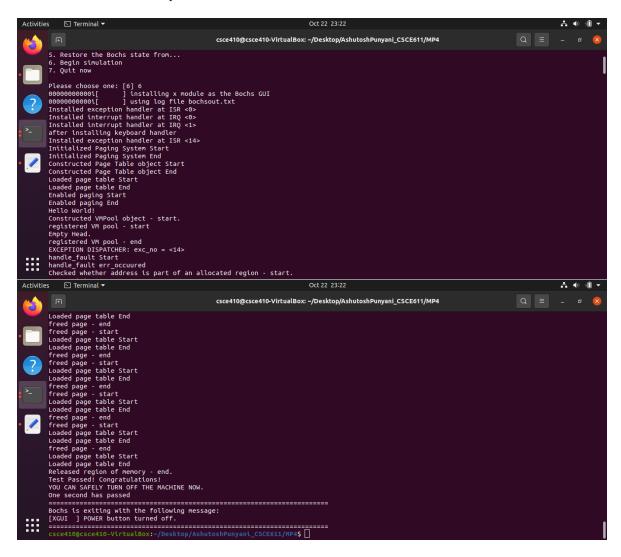


Figure 16: Testing (with commented _TEST_PAGE_TABLE_ define)

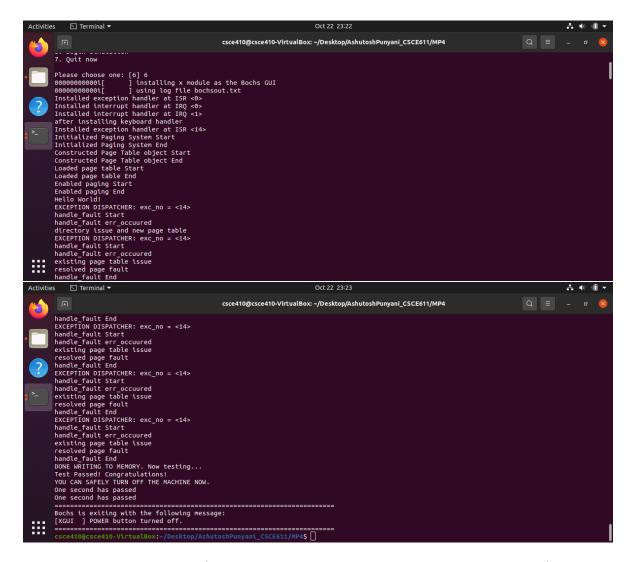


Figure 17: Testing (without commented _TEST_PAGE_TABLE_ define)