J8. Input: - x1, x2, ..., XN E R<sup>n</sup>,

Initial liet of K eluctor representatives

Z1, Zz, ..., ZK.

Output: - Cluster assignment c1, cz, ..., CN.

(a) cluster assignment based on cluster representatives. for each x; (i=1,2,...,N), we need to find ||xi-Z1||2, ||xi-Z2||2,..., ||xi-ZK||K

For computing  $||x_i - z_i||_2$ , no of subtractions = n, norm  $y \ge 2n$ , and there are k such computations. This amounts to (3nk).

Then, we also need to find the minimum among all K, that involves k computations.

So, cluster assignment for one x; takes (3n+1) K operations, and we have N such elements. So, total computational complexity = (3n+1) KN.

Ignoring the +1 and the constant factor of 3, we can say this is O(nKN), considering the by-0 notation.

(b) update cluster representatives. For all j = 1, 2, ..., Kdo  $\forall z_j = \frac{1}{|G_j|} \sum_{i \in G_j} x_i$ 

|     | The number of additions will be around N-K,   |
|-----|---|
|     | The number of additions will be around N-K, but as K < N, we can say no of additions N                        |
|     | Each addition requires n operations.  |
|     | Each addition requires no operations.  So, computations for addition = nN (n-length vector fakes no peration) |
|     | After that, number of divisions = K*n= kn   |
|     | So, total computational complexity for this step = $[nN+kn] = [O(n(N+k))]$                                    |
| (L) | Combining step I and 2, we can say that   |
|     | complianty = nKN+ nN+ kn  |
|     | As KCCN, and considering big-0 notation,  |
|     | me can wrute,   |
|     | complexity = 0 (nKN).   |
|     | So, in 10 iterations, no of computations of [10 nKN].   |
|     | Homener, if we want a more exact and  |
|     | jugorous bound, me can calculate that   |
|     | (3n+1) KN + (N-K) n + Kn x10  |
|     | etet! no. of additions ) no. of divisions   |
|     | = (3nkN+KN+ Nn)x10,   |