## STREAMING DATA ANALYSIS



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# STREAMING DATA ANALYSIS

- Introduction to the problem domain
- A few key data structures
- Implementation Considerations

Examples?

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Sensor data

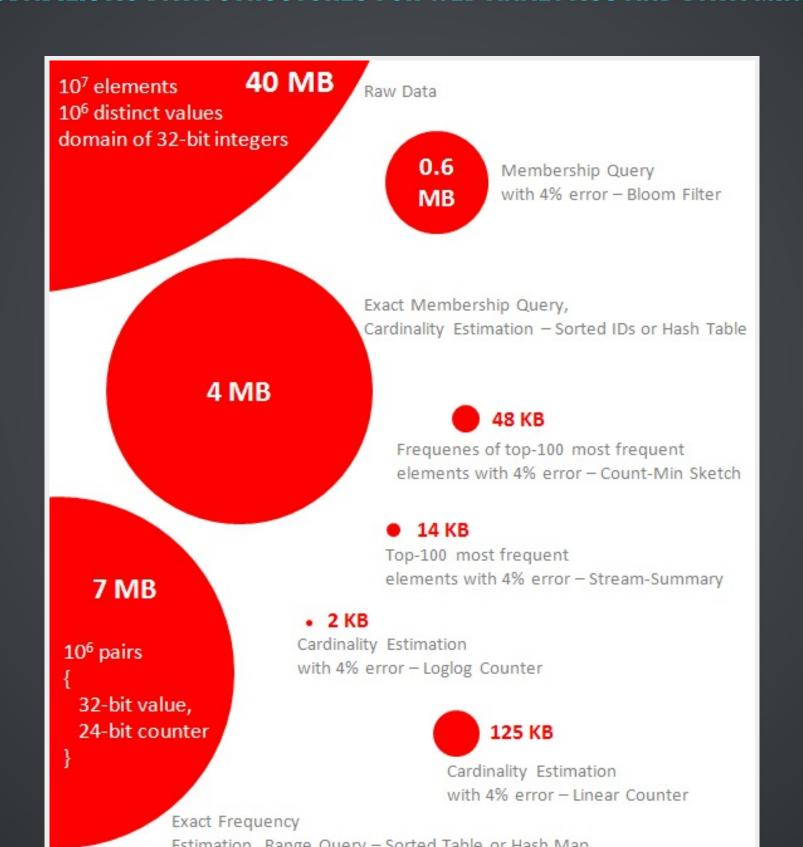
Examples?

- Sensor data
- Image data

Examples?

- Sensor data
- Image data
- Internet and Web traffic

#### PROBABILISTIC DATA STRUCTURES FOR WEB ANALYTICS AND DATA MINING



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- Many involve use of sketching



Set Membership

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Bloom Filter

Set Membership

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Cardinality Estimation

- Set MembershipBloom Filter
- Cardinality Estimation
  LogLog, HyperLogLog

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- Frequency Estimation

Set MembershipBloom Filter

Cardinality Estimation
 LogLog, HyperLogLog

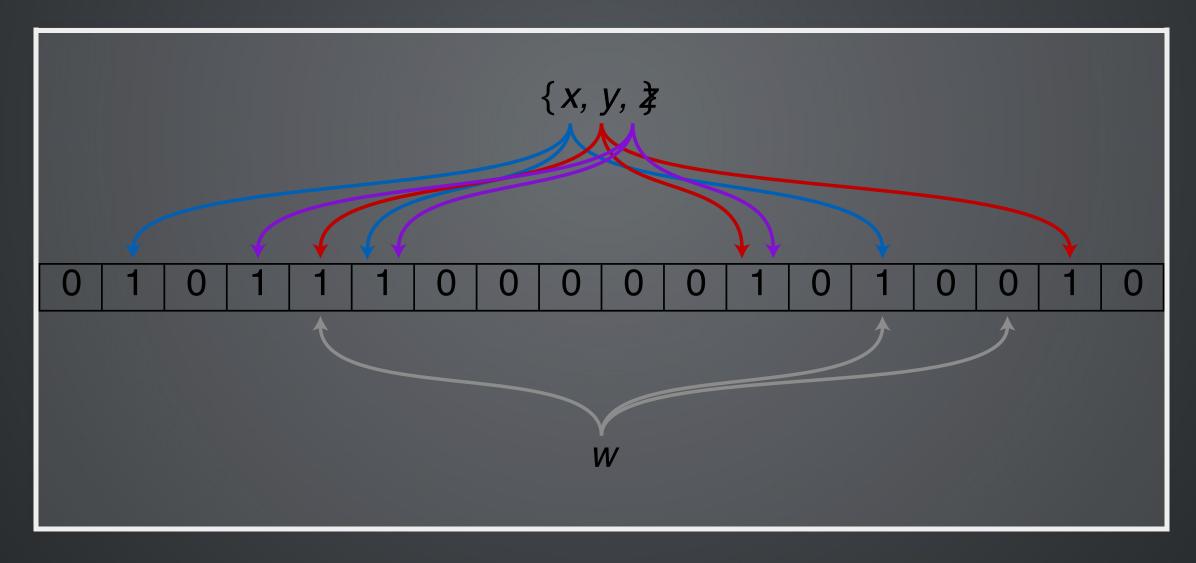
Frequency Estimation

Count Min Sketch, CountMean Min Sketch

- Set MembershipBloom Filter
- Cardinality Estimation
  LogLog, HyperLogLog
- Frequency Estimation
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- Heavy Hitters (Top-K)

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  LogLog, HyperLogLog
- Frequency Estimation
  Count Min Sketch, CountMean Min Sketch
- Heavy Hitters (Top-K)
  Count Min Sketch, Stream Summary

- Used to test whether an element is a member of a set
- False positive matches are possible, but false negatives are not



Algorithm

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- Define k different hash functions
- To add an element, feed it to each of the k hash functions to get k array positions
- Set the bits at all these positions to 1
- To query, feed it to each of the k hash functions to get k array positions. If any of the bits at these positions is 0, the element is definitely not in the set

Bloom Filter Demo

 Hash each element in the data set and represent as a binary string

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- Expect that about one half of strings will start with 1, one quarter will start with 01, and so on
- Denote the number of leading zeros as a rank
- If the maximum number of leading zeros observed is n, an estimate for the number of distinct elements in the set is  $2^n$

LogLog & HyperLogLog

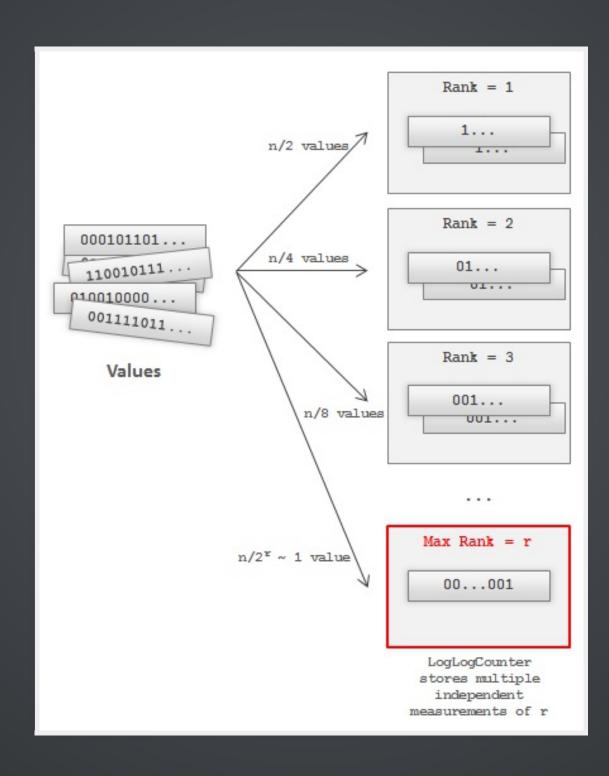
#### LogLog & HyperLogLog

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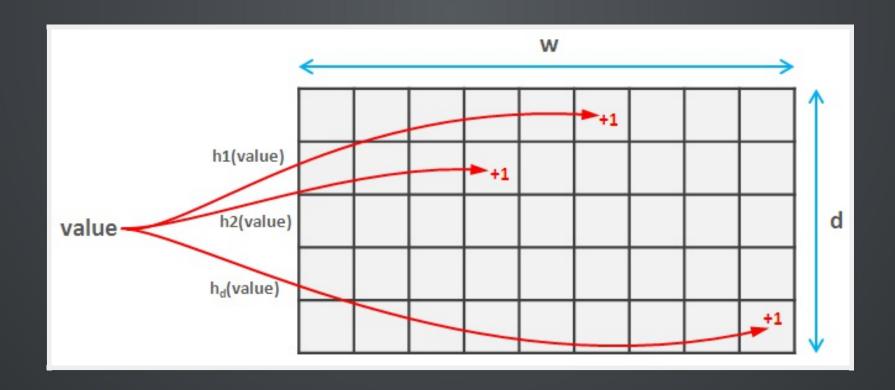
HLL is able to estimate cardinalities of  $> 10^9$  with a typical accuracy of 2%, using 1.5kB of memory



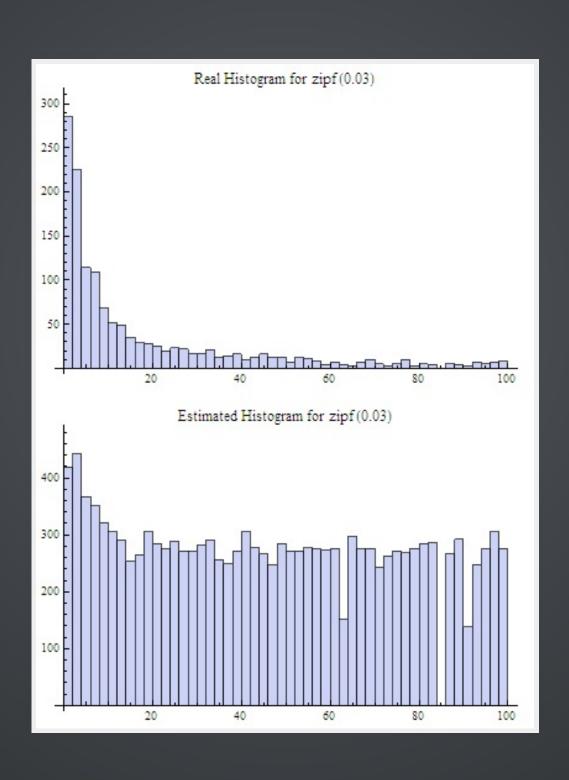
# COUNT MIN SKETCH

$$\epsilon \leq \frac{2n}{W}$$

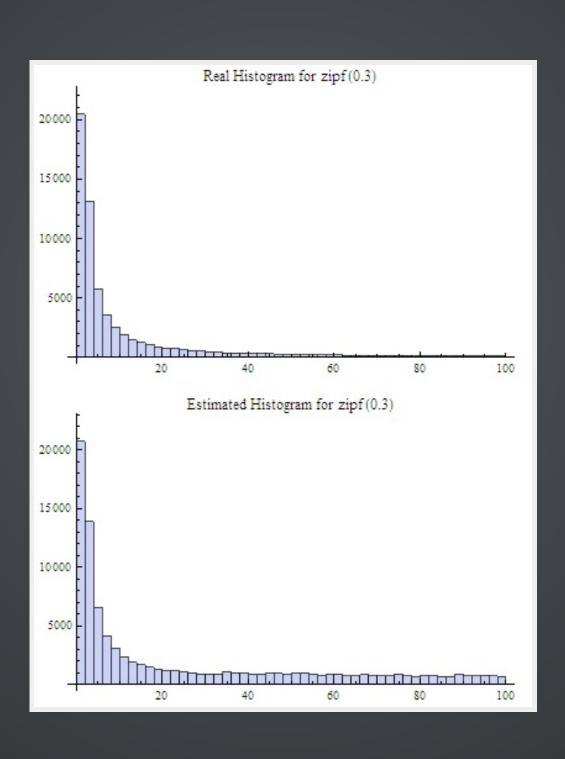
$$\delta = 1 - \left(\frac{1}{2}\right)^d$$



# COUNT MIN SKETCH

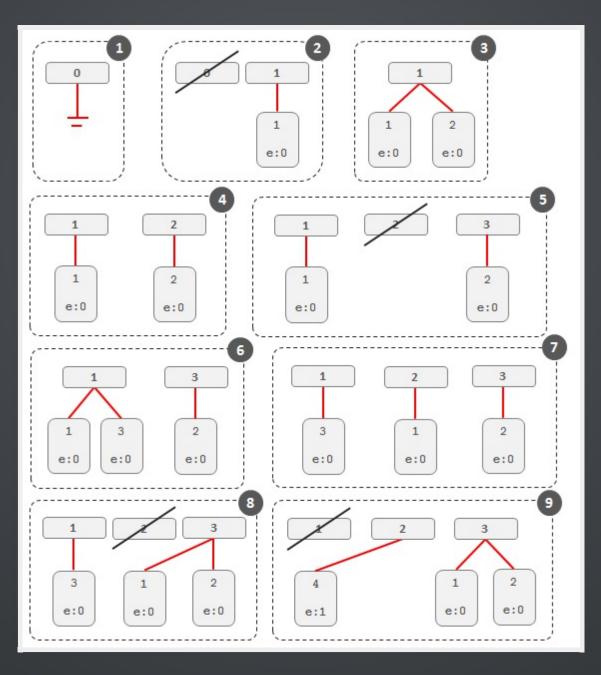


# COUNT MIN SKETCH

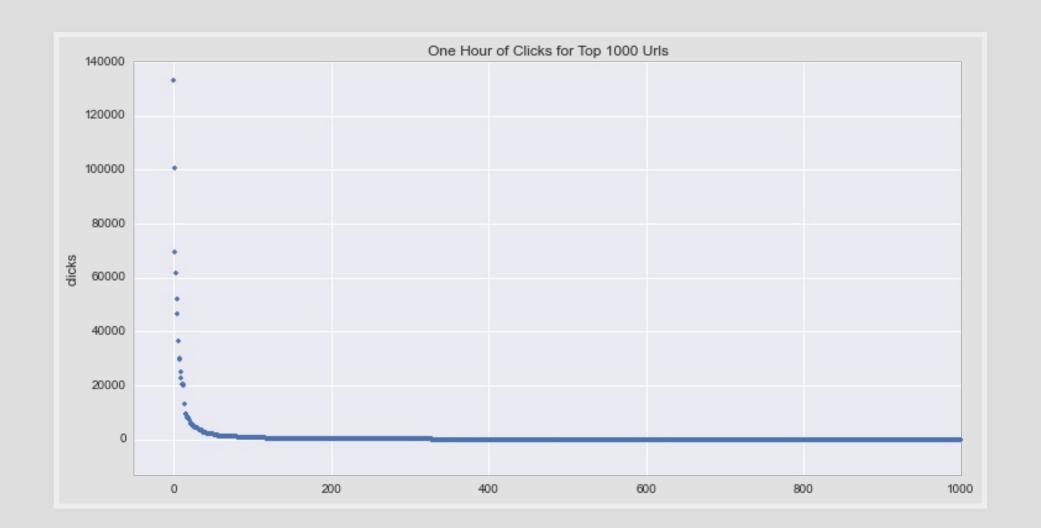


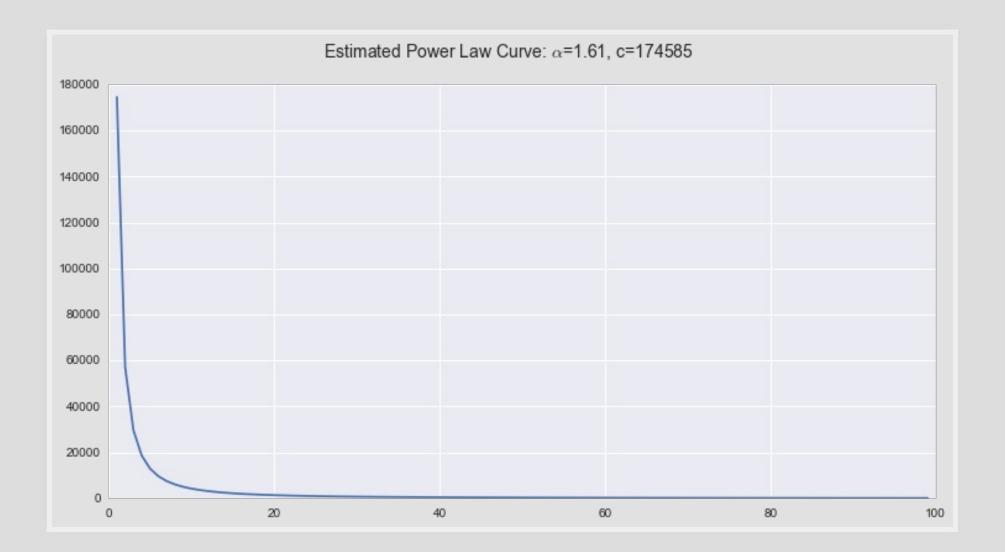
# STREAM SUMMARY

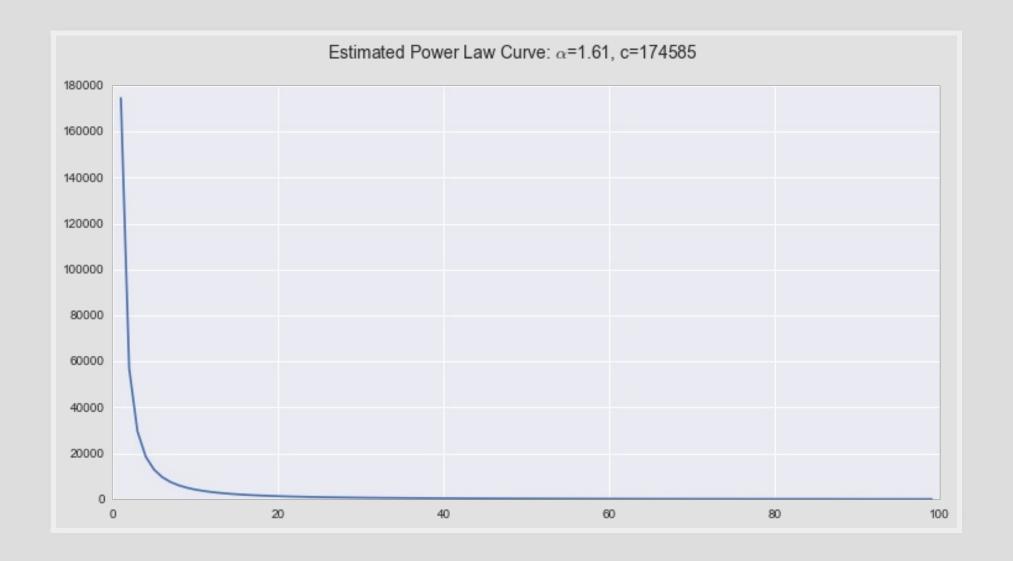
input stream: {1,2,2,3,1,1,4}



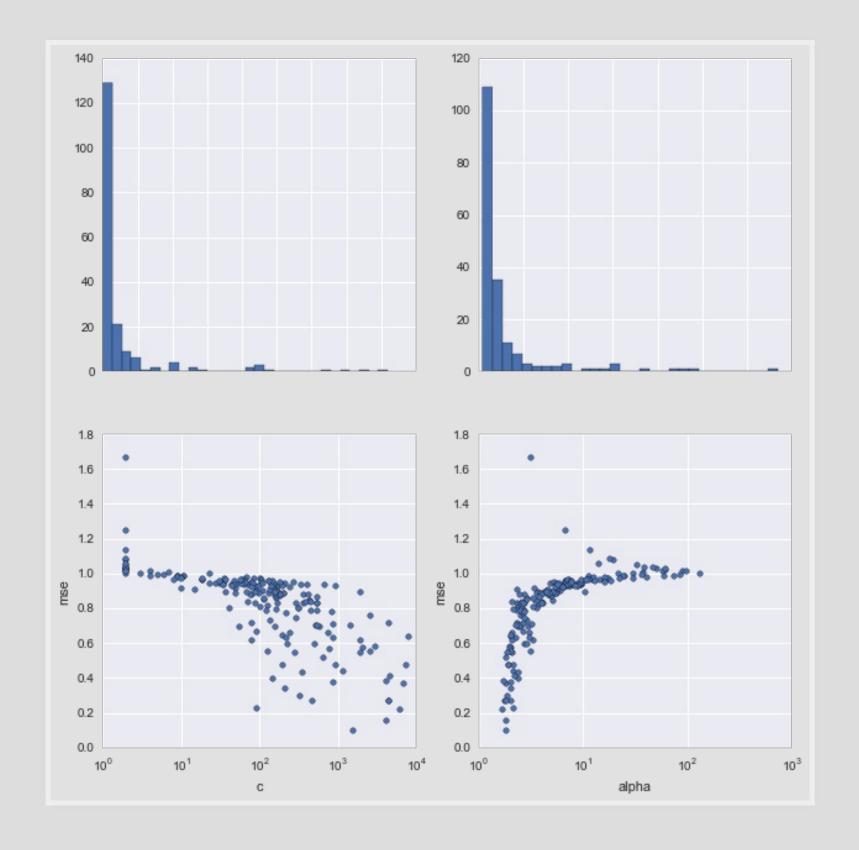
# IMPLEMENTATION CONSIDERATIONS







Important that  $\alpha > 1$ 



#### **FURTHER READING**

- Probabilistic Data Structures for Web Analytics and Data Mining, Ilya Katsov
- Efficient Computation of Frequent and Top-K Elements in Data Streams, A. Metwally, D. Agrawal, A.E. Abbadi.
- HyperLogLog: the analysis of a near-optimal cardinality estimation algorithm, P. Flayjolet, E.Fusy, O. Gandouet, F. Meunier.
- An Improved Data Stream Summary: The Count-Min Sketch and its Applications, . Cormode, S. Muthukrishnan.
- A Statistical Analysis of Probabilistic Counting Algorithms,
  P. Clifford, I. Cosma.
- Mining Massive Data Sets, Chapter 4Leskovec, Rajaraman, Ullman
- Stream-lib: Java library with implementations of many of