

(A Constituent College of Somaiya Vidyavihar University) **Department of Computer Engineering** 



Batch: B2 Roll No.: 16010124107

Experiment / assignment / tutorial No. 03

Grade: AA / AB / BB / BC / CC / CD /DD

Signature of the Staff In-charge with date

**TITLE:** To study and implement Restoring method of division

**AIM**: The basis of algorithm is based on paper and pencil approach and the operation involves repetitive shifting with addition and subtraction. So the main aim is to depict the usual process in the form of an algorithm.

**Expected OUTCOME of Experiment: (Mention CO /CO's attained here)** 

#### Books/ Journals/ Websites referred:

- **1.** Carl Hamacher, Zvonko Vranesic and Safwat Zaky, "Computer Organization", Fifth Edition, TataMcGraw-Hill.
- **2.** William Stallings, "Computer Organization and Architecture: Designing for Performance", Eighth Edition, Pearson.
- **3**. Dr. M. Usha, T. S. Srikanth, "Computer System Architecture and Organization", First Edition, Wiley-India.

### **Pre Lab/ Prior Concepts:**

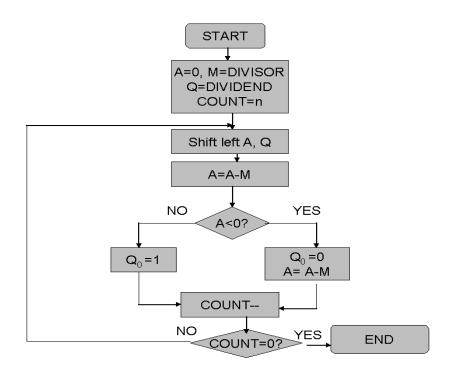
The Restoring algorithm works with any combination of positive and negative numbers

#### Flowchart for Restoring of Division:









### **Design Steps:**

- 1. Start
- 2. Initialize A=0, M=Divisor, Q=Dividend and count=n (no of bits)
- 3. Left shift A, Q
- 4. If MSB of A and M are same
- 5. Then A=A-M
- 6. Else A=A+M
- 7. If MSB of previous A and present A are same
- 8.  $Q_0=0$  & store present A
- 9. Else  $Q_0=0$  & restore previous A
- 10. Decrement count.
- 11. If count=0 go to 11
- 12. Else go to 3
- 13. STOP

### Example: - (Handwritten solved problem needs to be uploaded):-







9) 1	13		A=A+N
7) n	A	9	
4	0000	0001	SL
	0000	001-	A -= M
	1101	0010	<0? yes
	0000	0010	A=A+M
3	0000	010-	31
	1101	010 -	A=A-M
	*	0100	ACO? yes.
	0000	0100	ASA+M
2	0000	(00 -	SL
	1161	100 0	A-A-M
	0000	(000	ALO ? yes.
	0000	(000.	A=A+M
	00001	000 ±	A=A-M
	0001	000 -	
	1110	0000	A <0? yes
6	000 17 Kem	0000	A 2A +M

### CODE:-

#include <iostream>

#include <bitset>

using namespace std;

const int n = 8;







```
void leftShift(bitset<n>& A, bitset<n>& Q) {
  bool msbQ = Q[n - 1];
  A <<= 1;
  A[0] = msbQ;
  Q <<= 1;
}
void printState(int step, bitset<n> A, bitset<n> Q) {
  cout << "Step " << step << ": A = " << A << ", Q = " << Q << endl;
}
int main() {
  int dividend, divisor;
  cout << "Enter Dividend: ";
  cin >> dividend;
  cout << "Enter Divisor: ";</pre>
  cin >> divisor;
  bitset<n> A(0);
  bitset<n> Q(dividend);
  bitset<n> M(divisor);
  int count = n;
```







```
for (int step = 1; count > 0; --count, ++step) {
  leftShift(A, Q);
  int tempA = (int)(A.to_ulong()) - (int)(M.to_ulong());
  if (tempA < 0) {
     Q[0] = 0;
     tempA += M.to ulong();
  } else {
     Q[0] = 1;
  }
  A = bitset < n > (tempA);
  printState(step, A, Q);
}
cout << "Final Quotient (Q): " << Q << " = " << Q.to_ulong() << endl;
cout << "Final Remainder (A): " << A << " = " << A.to ulong() << endl;
return 0;
```

}







```
Enter positive Dividend: 4
Enter positive Divisor: 2
Step 1: A = 0000, Q = 1000
Step 2: A = 0001, Q = 0000
Step 3: A = 0000, Q = 0001
Step 4: A = 0000, Q = 0010
Final Quotient (Q): 0010 = 2
Final Remainder (A): 0000 = 0

Process returned 0 (0x0) execution time : 2.026 s
Press any key to continue.
```

### **Conclusion:-**

Restoring division is a method for performing binary division, commonly used in computer architecture. It involves repeatedly subtracting the divisor from the dividend and restoring the remainder if the result is negative. This process continues until the quotient and remainder are obtained.

### **Post Lab Descriptive Questions**

- 1. What are the advantages of restoring division over non restoring division?
  - The restoring division algorithm is conceptually simpler
  - Restoring division ensures that the remainder is always non-negative
  - Restoring division avoids the issue of sign correction at the end of the division process

**Date:** 01/08/2025