

Name: Ashwera Hasan
Experiment 4
Batch: B2 Roll: 16010124107

TITLE : To study and implement Non Restoring method of division

AIM : The basis of algorithm is based on paper and pencil approach and the operation involve repetitive shifting with addition and subtraction. So the main aim is to depict the usual process in the form of an algorithm.

Expected OUTCOME of Experiment: (Mention CO/CO's attained here)

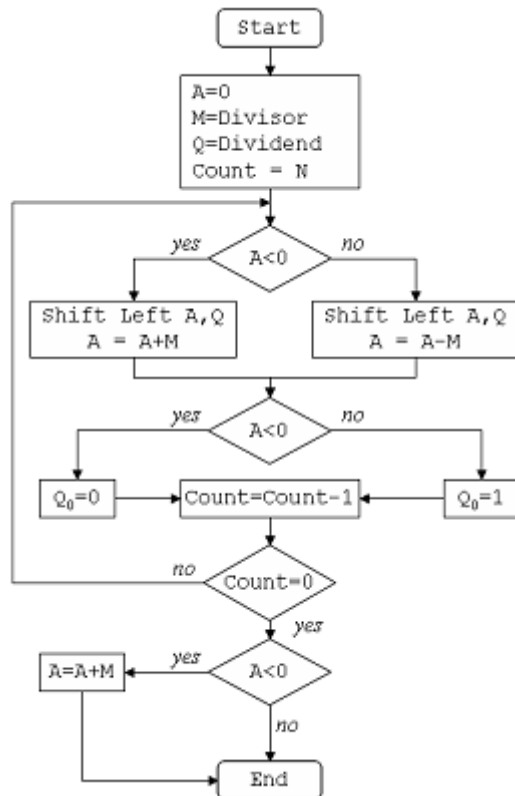
Books/ Journals/ Websites referred:

1. Carl Hamacher, Zvonko Vranesic and Safwat Zaky, "Computer Organization", Fifth Edition, TataMcGraw-Hill.
2. William Stallings, "Computer Organization and Architecture: Designing for Performance", Eighth Edition, Pearson.
3. Dr. M. Usha, T. S. Srikanth, "Computer System Architecture and Organization", First Edition, Wiley-India.

Pre Lab/ Prior Concepts:

The Non Restoring algorithm works with any combination of positive and negative numbers.

Flowchart for Non Restoring of Division(Students need to draw):-



Algorithm:

1. Start
2. Set $A=0$, $M=\text{divisor}$, and $Q = \text{dividend}$. Convert to binary. Set $N=\text{number of bits}+1$. Set $\text{count}=n+1$
3. Run this till $\text{count}=0$.
4. Shift A and Q to the left
5. If MSB of A is 1, do $A = A + M$
6. else, do $A - M = A$
7. after this operation, if MSB of A is 1, set the last bit of $Q=1$ else set 0
8. Reduce count.
9. Repeat steps 3 till 8.
10. If finally after $\text{count}=0$, MSB of A is 1, perform $A = A + M$
11. Quotient is in Q and Remainder is in A .
12. Stop.

Example: (Handwritten solved problem needs to uploaded):-

Q) $\frac{14}{3}$

$-M = -3 = 11101$
 $M = 3 = 00011$
 $Q = 14 = 01110$
 $A = 00000$

$(n+1) = 5$
 $n = 4$

A	B	n	Task
00000	01110	5	shift left
00000	1110 -		A & M have same sign: A < 0? NO.
11101	1110 0		A - M
11101	11100	4	A < 0, set = 0
11011	1100 -		SL
11110	11000		A < 0? yes. A + M
11110	11000	3	A < 0? yes set 0
11101	1000 -		SL
00000	1000 1		A < 0? A + M
00001	0001 -	2	A < 0? no
11110	0001 -		SL
			A < 0? A - M
11110	00010	1	A < 0? yes. 0
11100	0010 -		SL
11111	0010 -		A < 0 A + M
11111	00100	0	A < 0? 0.

$A = A + M = 00010 = 2$
 $Q = 00100 = 4$
 Rem = 00010 = 2.

CODE:-



```
nonrestoring.cpp
1  #include <iostream>
2  #include <iomanip>
3  using namespace std;
4
5  string toBinary(int num, int bits) {
6      string bin = "";
7      for (int i = bits - 1; i >= 0; --i)
8          bin += ((num >> i) & 1) ? '1' : '0';
9      return bin;
10 }
11
12 int main() {
13     int dividend, divisor;
14     cout << "Enter Dividend: ";
15     cin >> dividend;
16     cout << "Enter Divisor: ";
17     cin >> divisor;
18
19     if (dividend < 0 || divisor <= 0) {
20         cout << "Only positive integers supported.\n";
21         return 1;
22     }
23
24     int n = 4;
25     int A = 0;
26     int Q = dividend;
27     int M = divisor;
28
29     cout << "\nInitial values:\n";
30     cout << "n = " << n << ", A = " << A << ", Q = " << Q << ", M = " << M << "\n";
31 }
```

0 secs



```
nonrestoring.cpp > main()
int main() {
    int n = 16384;

    cout << "\nInitial values:\n";
    cout << "A = " << toBinary(A, n) << "\n";
    cout << "Q = " << toBinary(Q, n) << "\n";
    cout << "M = " << toBinary(M, n) << "\n";
    cout << "n = " << n << "\n\n";

    cout << left << setw(12) << "Step"
        << setw(10) << "A"
        << setw(10) << "Q"
        << setw(5) << "n"
        << "Task\n";
    cout << "-----\n";

    for (int step = 1; step <= n; ++step) {
        A = (A << 1) | ((Q >> (n - 1)) & 1); // MSB of Q to LSB of A
        Q = (Q << 1) & ((1 << n) - 1);      // Keep Q within n bits

        cout << setw(12) << step
            << setw(10) << toBinary(A, n)
            << setw(10) << toBinary(Q, n)
            << setw(5) << (n - step)
            << "Shift A and Q left\n";

        A = A - M;

        cout << setw(12) << ""
            << setw(10) << toBinary(A, n)

```

```
cc.cpp  submissionisallyouneed.cpp  nonrestoring.cpp X
nonrestoring.cpp > main()
12  int main() {
42      for (int step = 1; step <= n; ++step) {
53          cout << setw(12) << ""
54              << setw(10) << toBinary(A, n)
55              << setw(10) << toBinary(Q, n)
56              << setw(5) << ""
57              << "A = A - M\n";
58          if (A < 0) {
59              A = A + M;
60              Q = Q & (~1);
61              cout << setw(12) << ""
62                  << setw(10) << toBinary(A, n)
63                  << setw(10) << toBinary(Q, n)
64                  << setw(5) << ""
65                  << "A < 0, Restore A, Q0 = 0\n";
66          } else {
67              Q = Q | 1;
68              cout << setw(12) << ""
69                  << setw(10) << toBinary(A, n)
70                  << setw(10) << toBinary(Q, n)
71                  << setw(5) << ""
72                  << "A >= 0, Q0 = 1\n";
73          }
74      }
75  }
76
77  cout << "\nFinal Quotient (Q): " << toBinary(Q, n) << " (" << Q << ")\n";
78  cout << "Final Remainder (A): " << toBinary(A, n) << " (" << A << ")\n";
79
```

OUTPUT:-

```
g++ nonrestoring.cpp -o nonrestoring } ; if ($?) { .\nonrestoring
Enter Dividend: 14
Enter Divisor: 3

Initial values:
A = 0000
Q = 1110
M = 0011
n = 4
```

Step	A	Q	n	Task
1	0001	1100	3	Shift A and Q left
	1110	1100		A = A - M
	0001	1100		A < 0, Restore A, Q0 = 0
2	0011	1000	2	Shift A and Q left
	0000	1000		A = A - M
	0000	1001		A >= 0, Q0 = 1
3	0001	0010	1	Shift A and Q left
	1110	0010		A = A - M
	0001	0010		A < 0, Restore A, Q0 = 0
4	0010	0100	0	Shift A and Q left
	1111	0100		A = A - M
	0010	0100		A < 0, Restore A, Q0 = 0

```
Final Quotient (Q): 0100 (4)
Final Remainder (A): 0010 (2)
```

Conclusion:-

The non-restoring division is used for unsigned binary values that simplifies the procedure by eliminating the restoring phase. The non-restoring division is simpler and more effective than restoring division. It just employs addition and subtraction operations instead of restoring division, which requires extra steps to restore the original result after a failed subtraction.

Post Lab Descriptive Questions

Q. What are the advantages of non-restoring division over restoring division?

1. Non restoring division is faster and simpler
2. It is more effective
3. It avoids extra steps of restoring that used to obtain the previous A after a failed subtraction.
4. Requires less hardware in implementation.

Q. Solve $10/3$ using Non-Restoring algorithm for division operation?

10/3

M = 3 = 00011
Q = 10 = 01010
-M = 11100
1
11101

A	B	N	Task
00000	01010	5	shift
00000	1010 -		A < 0? no
11101	1010 -		A - M
			A < 0? yes Q = 0
11101	10100	4	shift
11011	0100 -		A < 0? A + M
11110	0100 -		A < 0? Q = 0
11110	01000	3	shift
11100	1000 -		A < 0? A + M
11111	1000 -		A < 0? Q = 0
11111	10000	2	shift
11111	0000 -		A < 0? A + M
00010	0000 -		A < 0? ^x Q = 1
00010	00001	1	shift
00100	0001 -		A < 0? ^x A - M
00001	0001 -		A < 0? ^b Q = 1
00001	00011	0	

Q = 00011 = 3
R = 00001 = 1
(3 x 3) + 1 = 10 ; verified ✓

Date:- 08.08.25