



Codd's 12 Rules for an RDBMS

Although most of us think that any database that supports SQL is automatically considered a relational database, this isn't always the case—at least not completely. In Chapter 1, I discussed the basics and foundations of relational theory, but no discussion on this subject would be complete without looking at the rules that were formulated in 1985 in a two-part article published by *Computerworld* magazine (“Is Your DBMS Really Relational?” and “Does Your DBMS Run By the Rules?” by E. F. Codd, *Computerworld*, October 14 and October 21, 1985). Many websites also outline these rules. These rules go beyond relational theory and defines more specific criteria that need to be met in an RDBMS, if it's to be truly relational.

It might seem like old news, but the same criteria can still be used today to measure how relational a database is. These rules are frequently brought up in conversations when discussing how well a particular database server is implemented. I present the rules in this appendix, along with brief comments as to whether SQL Server 2008 meets each of them, or otherwise. Relational theory has come a long way since these rules were first published, and “serious” theorists have enhanced and refined relational theory tremendously since then, as you'd expect. A good place for more serious learning is the website <http://www.dbdebunk.com>, run by C. J. Date and Fabian Pascal, or any of their books. If you want to see the debates on theory, the newsgroup `comp.databases.theory` is a truly interesting place. Of course, as the cover of this book states, my goal is practicality, with a foundation on theory, so I won't delve too deeply into theory here at all. I present these 12 rules simply to set a basis for what a relational database started out to be and largely what it is and what it should be even today.

All these rules are based upon what's sometimes referred to as the *foundation principle*, which states that for any system to be called a relational database management system, the relational capabilities must be able to manage it completely.

For the rest of this appendix, I'll treat SQL Server 2008 specifically as a relational database engine, not in any of the other configurations in which it might be used, such as a plain data store, a document storage device, or whatever other way you might use SQL Server as a storage engine.

Rule 1: The Information Rule

All information in the relational database is represented in exactly one and only one way—by values in tables.

This rule is an informal definition of a relational database and indicates that every piece of data that we permanently store in a database is located in a table.

In general, SQL Server fulfills this rule, because we cannot store any information in anything other than a table. We can't use the variables in this code to persist any data, and therefore they're scoped to a single batch.

Rule 2: Guaranteed Access Rule

Each and every datum (atomic value) is guaranteed to be logically accessible by resorting to a combination of table name, primary key value, and column name.

This rule stresses the importance of primary keys for locating data in the database. The table name locates the correct table, the column name finds the correct column, and the primary key value finds the row containing an individual data item of interest. In other words, each (atomic) piece of data is accessible by the combination of table name, primary key value, and column name. This rule exactly specifies how we access data using an access language such as Transact-SQL (T-SQL) in SQL Server.

Using SQL, we can search for the primary key value (which is guaranteed to be unique, based on relational theory, as long as it has been defined), and once we have the row, the data is accessed via the column name. We can also access data by any of the columns in the table, though we aren't always guaranteed to receive a single row back.

Rule 3: Systematic Treatment of NULL Values

NULL values (distinct from empty character string or a string of blank characters and distinct from zero or any other number) are supported in the fully relational RDBMS for representing missing information in a systematic way, independent of data type.

This rule requires that the RDBMS support a distinct NULL placeholder, regardless of datatype. NULLs are distinct from an empty character string or any other number, and they are always to be considered as unknown values.

NULLs must propagate through mathematic operations as well as string operations. $\text{NULL} + \langle \text{anything} \rangle = \text{NULL}$, the logic being that NULL means "unknown." If you add something known to something unknown, you still don't know what you have, so it's still unknown.

There are a few settings in SQL Server that can customize how NULLs are treated. Most of these settings exist because of some poor practices that were allowed in early versions of SQL Server:

- **ANSI_NULLS:** Determines how NULL comparisons are handled. When OFF, then $\text{NULL} = \text{NULL}$ is True for the comparison, and when ON (the default), $\text{NULL} = \text{NULL}$ returns UNKNOWN.
- **CONCAT_NULL_YIELDS_NULL:** When the **CONCAT_NULL_YIELDS_NULL** setting is set ON, NULLs are treated properly, such that $\text{NULL} + \text{'String Value'} = \text{NULL}$. If the **CONCAT_NULL_YIELDS_NULL** setting is OFF, which is allowed for backward compatibility with SQL Server, NULLs are treated in a nonstandard way such that $\text{NULL} + \text{'String Value'} = \text{'String Value'}$.

Rule 4: Dynamic Online Catalog Based on the Relational Model

The database description is represented at the logical level in the same way as ordinary data, so authorized users can apply the same relational language to its interrogation as they apply to regular data.

This rule requires that a relational database be self-describing. In other words, the database must contain certain system tables whose columns describe the structure of the database itself, or alternatively, the database description is contained in user-accessible tables.

This rule is becoming more of a reality in each new version of SQL Server, as with the implementation of the INFORMATION_SCHEMA system views. The INFORMATION_SCHEMA is a schema that has a set of views to look at much of the metadata for the tables, the relationships, the constraints, and even the code in the database.

Anything else you need to know can most likely be viewed in the system views (in the SYS schema). They're the system views that replaced the system tables in 2005 that we had used since the beginning of SQL Server time. These views are far easier to read and use, and most all the data is self-explanatory, rather than requiring bitwise operations on some columns to find the value.

Rule 5: Comprehensive Data Sublanguage Rule

A relational system may support several languages and various modes of terminal use. However, there must be at least one language whose statements are expressible, per some well-defined syntax, as character strings and whose ability to support all of the following is comprehensible: a. data definition b. view definition c. data manipulation (interactive and by program) d. integrity constraints e. authorization f. transaction boundaries (begin, commit, and rollback).

This rule mandates the existence of a relational database language, such as SQL, to manipulate data. SQL as such isn't specifically required. The language must be able to support all the central functions of a DBMS: creating a database, retrieving and entering data, implementing database security, and so on. T-SQL fulfils this function for SQL Server and carries out all the data definition and manipulation tasks required to access data.

SQL is a nonprocedural language, in that you don't specify "how" things happen, or even where. You simply ask a question of the relational server, and it does the work.

Rule 6: View Updating Rule

All views that are theoretically updateable are also updateable by the system.

This rule deals with views, which are virtual tables used to give various users of a database different views of its structure. It's one of the most challenging rules to implement in practice, and no commercial product fully satisfies it today.

A view is theoretically updateable as long as it's made up of columns that directly correspond to real table columns. In SQL Server, views are updateable as long as you don't update more than a single table in the statement; neither can you update a derived or constant field. SQL Server 2000 also implemented INSTEAD OF triggers that you can apply to a view (see Chapter 6). Hence, this rule can be technically fulfilled using INSTEAD OF triggers, but in what can be a less-than-straightforward manner. You need to take care when considering how to apply updates, especially if the view contains a GROUP BY clause and possibly aggregates.

Rule 7: High-Level Insert, Update, and Delete

The capability of handling a base relation or a derived relation as a single operand applies not only to the retrieval of data but also to the insertion, update, and deletion of data.

This rule stresses the set-oriented nature of a relational database. It requires that rows be treated as sets in insert, delete, and update operations. The rule is designed to prohibit implementations that support only row-at-a-time, navigational modification of the database. The SQL language covers this via the INSERT, UPDATE, and DELETE statements.

Even the CLR doesn't allow you to access the physical files where the data is stored, but BCP does kind of go around this. As always, you have to be extra careful when you use the low-level tools that can modify the data without going through the typical SQL syntax, because it can ignore the rules you have set up, introducing inconsistencies into your data.

Rule 8: Physical Data Independence

Application programs and terminal activities remain logically unimpaired whenever any changes are made in either storage representation or access methods.

Applications must still work using the same syntax, even when changes are made to the way in which the database internally implements data storage and access methods. This rule implies that the way the data is stored physically must be independent of the logical manner in which it's accessed. This is saying that users shouldn't be concerned about how the data is stored or how it's accessed. In fact, users of the data need only be able to get the basic definition of the data they need.

Other things that shouldn't affect the user's view of the data are as follows:

- *Adding indexes:* Indexes determine how the data is stored, yet the user, through SQL, will never know that indexes are being used.
- *Changing the filegroup of an object:* Just moving a table to a new filegroup will not affect the application. You access the data in the same way no matter where it is physically located.
- *Using partitioning:* Beyond moving entire tables around to different filegroups, you can move parts of a table around by using partitioning technologies to spread access around to different independent subsystems to enhance performance.
- *Modifying the storage engine:* From time to time, Microsoft has to modify how SQL Server operates (especially in major version upgrades). However, SQL statements must appear to access the data in the same manner as they did in any previous version, only (we hope) faster.

Microsoft has put a lot of work into this area, because SQL Server has a separate relational engine and storage engine, and OLE DB is used to pass data between the two. Further reading on this topic is available in SQL Server 2008 Books Online in the "Database Engine Components" topic or in *Inside Microsoft SQL Server 2005: The Storage Engine* by Kalen Delaney (Microsoft Press, 2006).

Rule 9: Logical Data Independence

Application programs and terminal activities remain logically unimpaired when information preserving changes of any kind that theoretically permit unimpairment are made to the base tables.

Along with rule 8, this rule insulates the user or application program from the low-level implementation of the database. Together, they specify that specific access or storage techniques used by the RDBMS—and even changes to the structure of the tables in the database—shouldn't affect the user's ability to work with the data.

In this way, if you add a column to a table and if tables are split in a manner that doesn't add or subtract columns, then the application programs that call the database should be unimpaired.

For example, say you have the table in Figure A-1.

baseTable	
baseTableId	
column1	
column2	

Figure A-1. Small sample table

Then, say you vertically break it up into two tables (see Figure A-2).

baseTableA	
baseTableId	
column1	

baseTableB	
baseTableId	
column2	

Figure A-2. Small sample table split into two tables

You then could create this view:

```
CREATE VIEW baseTable
AS
SELECT baseTableId, column1, column2
FROM   baseTableA
       JOIN baseTableB
         ON baseTableA.baseTableId = baseTableB.baseTableId
```

The user should be unaffected. If you were to implement `INSTEAD OF` triggers on the view that had the same number of columns with the same names, you could seamlessly meet the need to manage the view in the exact manner the table was managed. Note that the handling of identity columns can be tricky in views, because they require data to be entered, even when the data won't be used. See Chapter 6 for more details on creating `INSTEAD OF` triggers.

Of course, you cannot always make this rule work if columns or tables are removed from the system, but you can make the rule work if columns and data are simply added.

Tip Always access data from the RDBMS by name, and not by position or by using the `SELECT *` wildcard. The order of columns shouldn't make a difference to the application.

Rule 10: Integrity Independence

Integrity constraints specific to a particular relational database must be definable in the relational data sublanguage and storable in the catalog, not in the application programs.

The database must support a minimum of the following two integrity constraints:

- *Entity integrity*: No component of a primary key is allowed to have a NULL value.
- *Referential integrity*: For each distinct non-NULL foreign key value in a relational database, there must exist a matching primary key value from the same domain.

This rule says that the database language should support integrity constraints that restrict the data that can be entered into the database and the database modifications that can be made. In other words, the RDBMS must internally support the definition and enforcement of entity integrity (primary keys) and referential integrity (foreign keys).

It requires that the database be able to implement constraints to protect the data from invalid values and that careful database design is needed to ensure that referential integrity is achieved. SQL Server 2008 does a great job of providing the tools to make this rule a reality. We can protect our data from invalid values for most any possible case using constraints and triggers. Most of Chapter 6 was spent covering the methods that we can use in SQL Server to implement integrity independence.

Rule 11: Distribution Independence

The data manipulation sublanguage of a relational DBMS must enable application programs and terminal activities to remain logically unimpaired whether and whenever data are physically centralized or distributed.

This rule says that the database language must be able to manipulate data located on other computer systems. In essence, we should be able to split the data on the RDBMS out onto multiple physical systems without the user realizing it. SQL Server 2008 supports distributed transactions among SQL Server sources, as well as other types of sources using the Microsoft Distributed Transaction Coordinator service.

Another distribution-independence possibility is a group of database servers working together more or less as one. Database servers working together like this are considered to be *federated*. With every new SQL Server version, the notion of federated database servers seamlessly sharing the load is becoming a definite reality. More reading on this subject can be found in the SQL Server 2008 Books Online in the “Federated Database Servers” topic.

Rule 12: Non-Subversion Rule

If a relational system has or supports a low-level (single-record-at-a-time) language, that low-level language cannot be used to subvert or bypass the integrity rules or constraints expressed in the higher-level (multiple-records-at-a-time) relational language.

This rule requires that alternate methods of accessing the data are not able to bypass integrity constraints, which means that users can't violate the rules of the database in any way. For most SQL

Server 2008 applications, this rule is followed, because there are no methods of getting to the raw data and changing values other than by the methods prescribed by the database. However, SQL Server 2008 violates this rule in two places:

- *Bulk copy*: By default, you can use the bulk copy routines to insert data into the table directly and around the database server validations.
- *Disabling constraints and triggers*: There's syntax to disable constraints and triggers, thereby subverting this rule.

It's always good practice to make sure you use these two features carefully. They leave gaping holes in the integrity of your data, because they allow any values to be inserted in any column. Because you're expecting the data to be protected by the constraint you've applied, data value errors might occur in the programs that use the data, without revalidating it first.

Summary

Codd's 12 rules for relational databases can be used to explain much about how SQL Server operates today. These rules were a major step forward in determining whether a database vendor could call his system "relational" and presented stiff implementation challenges for database developers. Fifteen years on, even the implementation of the most complex of these rules is becoming achievable, though SQL Server (and other RDBMSs) still fall short of achieving all their objectives.



Scalar Datatype Reference

Choosing proper datatypes to match the domain to satisfy logical modeling is an important task. One datatype might be more efficient than another of a similar type. For example, you can store integer data in an integer datatype, a numeric datatype, a floating point datatype, a character type, or even a binary column, but these datatypes certainly aren't alike in implementation or performance.

In this appendix, I'll introduce you to all the intrinsic datatypes that Microsoft provides and discuss the situations where they're best used. The following is a list of the datatypes I'll cover. I'll discuss when to use them and in some cases why not to use them.

- *Precise numeric data*: Stores data with no loss of precision due to storage.
 - bit: Stores either 1, 0, or NULL. Used for Boolean-type columns.
 - tinyint: Non-negative values between 0 and 255.
 - smallint: Integers between -32,768 and 32,767.
 - int: Integers between -2,147,483,648 and 2,147,483,647 (-2^{31} to $2^{31} - 1$).
 - bigint: Integers between -9,223,372,036,854,775,808 and 9,223,372,036,854,775,807 (that is, -2^{63} to $2^{63} - 1$).
 - decimal: Values between $-10^{38} + 1$ through $10^{38} - 1$.
 - money: Values from -922,337,203,685,477.5808 through 922,337,203,685,477.5807.
 - smallmoney: Values from -214,748.3648 through +214,748.3647.
- *Approximate numeric data*: Stores approximations of numbers. Provides for a large range of values.
 - float (N): Values in the range from $-1.79E + 308$ through $1.79E + 308$.
 - real: Values in the range from $-3.40E + 38$ through $3.40E + 38$. real is a synonym for a float(24) datatype.
- *Date and time*: Stores date values, including time of day.
 - date: Date-only values from January 1, 0001, to December 31, 9999 (3 bytes).
 - time: Time of day-only values to 100 nanoseconds (3–5 bytes). Note that the range of this type is from 0:00 to 23:59:59 and some fraction of a second based on the precision you select.
 - smalldatetime: Dates from January 1, 1900, through June 6, 2079, with accuracy to 1 minute (4 bytes).
 - datetime: Dates from January 1, 1753, to December 31, 9999, with accuracy to ~3 milliseconds (stored in increments of .000, .003, or .007 seconds) (8 bytes).

- **datetime2**: Despite the hideous name, this type will store dates from January 1, 0001, to December 31, 9999, to 100-nanosecond accuracy (6–8 bytes). The accuracy is based on the precision you select.
- **datetimeoffset**: Same as **datetime2** but includes an offset from UTC time (8–10 bytes).
- **Character (or string) data**: Used to store textual data, such as names, descriptions, notes, and so on.
 - **char**: Fixed-length character data up to 8,000 characters long.
 - **varchar**: Variable-length character data up to 8,000 characters long.
 - **varchar(max)**: Large variable-length character data; maximum length of $2^{31} - 1$ (2,147,483,647) bytes, or 2GB.
 - **text**: Large text values; maximum length of $2^{31} - 1$ (2,147,483,647) bytes, or 2GB. (Note that this datatype is outdated and should be phased out in favor of the **varchar(max)** datatype.)
 - **nchar, nvarchar, ntext**: Unicode equivalents of **char**, **varchar**, and **text** (with the same deprecation warning for **ntext** as for **text**).
- **Binary data**: Data stored in bytes, rather than as human-readable values, for example, files or images.
 - **binary**: Fixed-length binary data up to 8,000 bytes long.
 - **varbinary**: Variable-length binary data up to 8,000 bytes long.
 - **varbinary(max)**: Large binary data; maximum length of $2^{31} - 1$ (2,147,483,647) bytes, or 2GB.
 - **image**: Large binary data; maximum length of $2^{31} - 1$ (2,147,483,647) bytes, or 2GB. (Note that this datatype is outdated and should be phased out for the **varbinary(max)** datatype.)
- **Other scalar datatypes**: Datatypes that don't fit into any other groups nicely but are still interesting.
 - **timestamp (or rowversion)**: Used for optimistic locking.
 - **uniqueidentifier**: Stores a globally unique identifier (GUID) value.
 - **cursor**: Datatype used to store a cursor reference in a variable. Cannot be used as a column in a table.
 - **table**: Used to hold a reference to a local temporary table. Cannot be used as a column in a table.
 - **sql_variant**: Stores data of most any datatype.
- **Not simply scalar**: For completeness, I will mention these types, but they will be covered on their own in Appendix C (which will be available as bonus downloadable content). These types are XML, **hierarchyid**, and the spatial types (geometry and geography).

Although you'll look at all these datatypes, this doesn't mean you'll have a need for all of them. Choosing a datatype needs to be a specific task to meet the needs of the client with the proper datatype. You could just store everything in unlimited-length character strings (this was how some systems worked in the old days), but this is clearly not optimal. From the list, you'll choose the best datatype, and if you cannot find one good enough, you can use the CLR and implement your own (I cover this in Chapter 5). The proper datatype choice is the first step in making sure the proper data is stored for a column.

Note I include information in each section about how the types are affected by using compression. This information refers to row-level compression only. For page-level compression information, see Chapter 9. Also note that compression is available only in the Enterprise Edition of SQL Server 2008.

Precise Numeric Data

You can store numerical data in many base datatypes. There are two different types of numeric data: precise and approximate. The differences are important and must be well understood by any architect who's building a system that stores readings, measurements, or other numeric data.

Precise values have no error in the way they're stored, from integer to floating point values, because they have a fixed number of digits before and after the decimal point (or *radix*). It might seem odd to need to say that, but as I'll discuss in the next major section, some datatypes are considered approximate in that they don't always store exactly what you expect them to store. However, they aren't as bad as they sound and are useful for scientific and other applications where the range of values varies greatly.

The precise numeric values include the `bit`, `int`, `bigint`, `smallint`, `tinyint`, `decimal`, and `money` datatypes (`money` and `smallmoney`). I'll break these down again into two additional subsections: whole numbers and fractional numbers. This is done so we can isolate some of the discussion down to the values that allow fractional parts to be stored, because quite a few mathematical "quirks" need to be understood surrounding using those datatypes. I'll mention a few of these quirks, most importantly with the `money` datatypes. However, when you do any math with computers, you must be careful how rounding is achieved and how this affects your results.

Integer Values

Whole numbers are, for the most part, integers stored using base-2 values. You can do bitwise operations on them, though generally it's frowned upon in SQL (think back to Chapter 4 in the First Normal Form sections, if you don't remember why). Math performed with integers is generally fast because the CPU can perform it directly using registers. I'll cover five integer sizes: `bit`, `tinyint`, `smallint`, `int`, and `bigint`.

bit

Domain: 0, 1, or NULL.

Storage: A `bit` column requires 1 byte of storage per eight instances in a table. Hence, having eight `bit` columns will cause your table to be no larger than if your table had only a single `bit` column.

Discussion:

You use `bit` values as a kind of imitation Boolean value. A `bit` isn't a Boolean value, in that it has values 0 and 1, not `True` and `False`. This is a minor distinction but one that needs to be made. You cannot execute code such as this:

```
IF (bitValue) DO SOMETHING
```

A better term than a Boolean is a *flag*. A value of 1 means the flag has been set (such as a value that tells us that a customer does want e-mail promotions). Many programmers like to use character values 'yes' or 'no' for this, because this can be easier for viewing, but it can be harder to program with using built-in programming methods. In fact, the use of the bit datatype as a Boolean value has occurred primarily because many programming languages usually use 0 for False and nonzero for True (some use 1 or -1 explicitly).

You can index a bit column, but usually it isn't of any value only to index it. Having only two distinct values in an index (technically three with NULL) makes for a poor index. (See Chapter 9 for more information about indexes. You may be able to use a filtered index to make some indexes on bit columns useful.) Clearly, a bit value most often should be indexed in conjunction with other columns.

Row Compression Effect:

This will require 4 bits because of the metadata overhead of compression.

Tip There's always a ton of discussion on the newsgroups about using the bit datatype. It isn't a standard datatype, but in my opinion it's no different in usage than an integer that has been constrained to the values 0, 1, or NULL.

It's often asked why we don't have a Boolean datatype. This is largely because of the idea that datatypes need to support NULL in RDBMSs, and a Boolean datatype would have to support UNKNOWN and NULL, resulting in four valued logic tables that are difficult to contemplate (without taking a long nap) and hard to deal with. So, we have what we have, and it works well enough.

tinyint

Domain: Non-negative whole numbers from 0 through 255.

Storage: 1 byte.

Discussion:

tinyints are used to store small non-negative integer values. When using a single byte for storage, if the values you'll be dealing with are guaranteed always to be in this range, a tinyint is perfect. A great use for this is for the primary key of a domain table that can be guaranteed to have only a couple values. The tinyint datatype is especially useful in a data warehouse to keep the surrogate keys small. However, you have to make sure that there will never be more than 256 values, so unless the need for performance is incredibly great (such as if the key will migrate to tables with billions and billions of rows), it's best to use a larger datatype.

Row Compression Effect:

No effect, because 1 byte is the minimum for the integer types other than bit.

smallint

Domain: Whole numbers from -32,768 through 32,767 (or -2^{15} through $2^{15} - 1$).

Storage: 2 bytes.

Discussion:

If you can be guaranteed to need values only in this range, the `smallint` can be a useful type. It requires 2 bytes of storage.

Just as before with `tinyint`, it's often a bad idea to use a `smallint` for a primary key. Uniformity (just using `int`) makes your database code more consistent. This might seem like a small point, but in most average systems it's much easier to code when you automatically know what the datatype is.

One use of a `smallint` that crops up from time to time is as a Boolean. This is because, in earlier versions of Visual Basic, 0 equals False and -1 equals True (technically, VB would treat any nonzero value as True, but it used -1 as a False). Storing data in this manner is not only a tremendous waste of space—2 bytes versus potentially 1/8th of a byte for a bit, or even a single byte for a char(1)—'Y' or 'N'. It's also confusing to all the other SQL Server programmers. ODBC and OLE DB drivers do this translation for you, but even if they didn't, it's worth the time to write a method or a function in VB to translate True to a value of 1.

Row Compression Effect:

The value will be stored in the smallest number of bytes required to represent the value. For example, if the value is 10, it would fit in a single byte; then it would use 1 byte, and so forth, up to 2 bytes.

int

Domain: Whole numbers from -2,147,483,648 to 2,147,483,647 (that is, -2^{31} to $2^{31} - 1$).

Storage: 4 bytes.

Discussion:

The integer datatype is frequently employed as a primary key for tables, because it's small (it requires 4 bytes of storage) and efficient to store and retrieve.

One downfall of the `int` datatype is that it doesn't include an unsigned version, which for a 32-bit version could store non-negative values from 0 to 4,294,967,296 (2^{32}). Because most primary key values start out at 1, this would give you more than 2 billion extra values for a primary key value without having to involve negative numbers that can be confusing to the user. This might seem unnecessary, but systems that have billions of rows are becoming more and more common.

An application where the storage of the `int` column plays an important part is the storage of IP addresses as integers. An IP address is simply a 32-bit integer broken down into four octets. For example, if you had an IP address of 234.23.45.123, you would take $(234 * 2^3) + (23 * 2^2) + (45 * 2^1) + (123 * 2^0)$. This value fits nicely into an unsigned 32-bit integer, but not into a signed one. However, the 64-bit integer (`bigint`, which is covered next) in SQL Server covers the current IP address standard nicely but requires twice as much storage. Of course, `bigint` will fall down in the same manner when we get to IPv6 (the forthcoming Internet addressing protocol), because it uses a full 64-bit unsigned integer.

Row Compression Effect:

The value will be stored in the smallest number of bytes required to represent the value. For example, if the value is 10, it would fit in a single byte; then it would use 1 byte, and so forth, up to 4 bytes.

bigint

Domain: Whole numbers from -9,223,372,036,854,775,808 to 9,223,372,036,854,775,807 (that is, -2^{63} to $2^{63} - 1$).

Storage: 8 bytes.

Discussion:

One of the common reasons to use the 64-bit datatype is as a primary key for tables where you'll have more than 2 billion rows, or if the situation directly dictates it, such as the IP address situation I previously discussed. Of course, there are some companies where a billion isn't really a very large number of things to store or count, so using a `bigint` will be commonplace to them. As usual, the important thing is to size your utilization of any type to the situation, not using too small or even too large of a type than is necessary.

Row Compression Effect:

The value will be stored in the smallest number of bytes required to represent the value. For example, if the value is 10, it would fit in a single byte; then it would use 1 byte, and so forth, up to 4 bytes.

Decimal Values

The decimal datatype is precise, in that whatever value you store, you can always retrieve it from the table. However, when you must store fractional values in precise datatypes, you pay a performance and storage cost in the way they're stored and dealt with. The reason for this is that you have to perform math with the precise decimal values using SQL Server engine code. On the other hand, math with IEEE floating point values (the `float` and `real` datatypes) can use the floating point unit (FPU), which is part of the core processor in all modern computers. This isn't to say that the decimal type is slow, *per se*, but if you're dealing with data that doesn't require the perfect precision of the decimal type, use the `float` datatype. I'll discuss the `float` and `real` datatypes more in the "Approximate Numeric Data" section.

decimal (or Numeric)

Domain: All numeric data (including fractional parts) between $-10^{38} + 1$ through $10^{38} - 1$.

Storage: Based on precision (the number of significant digits): 1–9 digits, 5 bytes; 10–19 digits, 9 bytes; 20–28 digits, 13 bytes; and 29–38 digits, 17 bytes.

Discussion:

The decimal datatype is a precise datatype because it's stored in a manner that's like character data (as if the data had only 12 characters, 0 to 9 and the minus and decimal point symbols). The way it's stored prevents the kind of imprecision you'll see with the `float` and `real` datatypes a bit later. However, `decimal` does incur an additional cost in getting and doing math on the values, because there's no hardware to do the mathematics.

To specify a decimal number, you need to define the precision and the scale:

- *Precision* is the total number of significant digits in the number. For example, 10 would need a precision of 2, and 43.00000004 would need a precision of 10. The precision may be as small as 1 or as large as 38.
- *Scale* is the possible number of significant digits to the right of the decimal point. Reusing the previous example, 10 would have a scale of 0, and 43.00000004 would need 8.

Numeric datatypes are bound by this precision and scale to define how large the data is. For example, take the following declaration of a numeric variable:

```
DECLARE @testvar decimal(3,1)
```

This allows you to enter any numeric values greater than -99.94 and less than 99.94. Entering 99.949999 works, but entering 99.95 doesn't, because it's rounded up to 100.0, which can't be displayed by decimal(3,1). Take the following, for example:

```
SELECT @testvar = -10.155555555
SELECT @testvar
```

This returns the following result:

```
-----
-10.2
```

This rounding behavior is both a blessing and a curse. You must be careful when butting up to the edge of the datatype's allowable values. Note that a setting—`SET NUMERIC_ROUNDABORT ON`—causes an error to be generated when a loss of precision would occur from an implicit data conversion. That's kind of like what happens when you try to put too many characters into a character value.

Take the following code:

```
SET NUMERIC_ROUNDABORT ON
DECLARE @testvar decimal(3,1)
SELECT @testvar = -10.155555555
```

This causes the following error:

```
Msg 8115, Level 16, State 7, Line 5
Arithmetic overflow error converting numeric to data type numeric.
```

`SET NUMERIC_ROUNDABORT` can be quite dangerous to use and might throw off applications using SQL Server if set to `ON`. However, if you need to prevent implicit round-off due to system constraints, it's there.

As far as usage is concerned, you should generally use the `decimal` datatype as sparingly as possible, and I don't mean this negatively. There's nothing wrong with the type at all, but it does take that little bit more processing than integers or real data, and hence there's a performance hit. You should use it when you have specific values that you want to store where you can't accept any loss of precision. Again, I'll deal with the topic of loss of precision in more detail in the section "Approximate Numeric Data." The `decimal` type is commonly used as a replacement for the money type, because it has certain round-off issues that `decimal` does not.

Row Compression Effect:

The value will be stored in the smallest number of bytes that are necessary to provide the precision necessary, plus 2 bytes overhead per row. For example, if you are storing the value of 2 in a `numeric(28,2)` column, it needn't use all the possible space; it can use the space of a `numeric(3,2)`, plus the 2 bytes overhead.

Money Types

There are two intrinsic datatypes that are for storing monetary values. Both are based on integer types, with a fixed four decimal places. These types are as follows:

- `money`
 - *Domain*: -922,337,203,685,477.5808 to 922,337,203,685,477.5807
 - *Storage*: 8 bytes
- `smallmoney`
 - *Domain*: -214,748.3648 to 214,748.3647
 - *Storage*: 4 bytes

The money datatypes are generally considered a poor choice of datatype, even for storing monetary values, because they have a few inconsistencies that can cause a good deal of confusion. First, you can specify units, such as \$ or £, but the units are of no real value. For example:

```
CREATE TABLE dbo.testMoney
(
    moneyValue money
)
go

INSERT INTO dbo.testMoney
VALUES ($100)
INSERT INTO dbo.testMoney
VALUES (100)
INSERT INTO dbo.testMoney
VALUES (£100)
GO
SELECT * FROM dbo.testMoney
```

The query at the end of this code example returns the following results:

```
moneyValue
-----
100.00
100.00
100.00
```

The second problem is that the money datatypes have well-known rounding issues with math. I mentioned that these types are based on integers (the range for `smallmoney` is -214,748.3648 to 214,748.3647, and the range for an integer is 2,147,483,648 to 2,147,483,647). Unfortunately, as I will demonstrate, intermediate results are stored in the same types, causing unexpected rounding errors. For example:

```
DECLARE @money1 money, @money2 money

SET @money1 = 1.00
SET @money2 = 800.00
SELECT cast(@money1/@money2 as money)
```

This returns the following result:

```
-----
0.0012
```

However, try the following code:

```
DECLARE @decimal1 decimal(19,4), @decimal2 decimal(19,4)
SET      @decimal1 = 1.00
SET      @decimal2 = 800.00
SELECT   cast(@decimal1/@decimal2 as decimal(19,4))
```

It returns the following result:

```
-----
0.0013
```

Why? Because `money` uses only four decimal places for intermediate results, where `decimal` uses a much larger precision:

```
SELECT @money1/@money2
SELECT @decimal1/@decimal2
```

This code returns the following results:

```
-----
0.0012
```

```
-----
0.0012500000000000000
```

That's why there are round-off issues. The common consensus among database architects is to avoid the `money` datatype and use a numeric type instead, because of the following reasons:

- It gives the answers to math problems in the natural manner that's expected.
- It has no built-in units to confuse matters.

Even in the previous version of SQL Server, the following statement was included in the monetary data section: "If a greater number of decimal places are required, use the `decimal` datatype instead." Using a `decimal` type instead gives you the precision needed. To replicate the range for `money`, use `DECIMAL (19,4)`, or for `smallmoney`, use `DECIMAL (10,4)`. However, you needn't use such large values if you don't need them. Otherwise, if you happen to be calculating the national debt or my yearly gadget allowance, you might need to use a larger value.

Row Compression Effect:

The `money` types are simply integer types with their decimal places shifted. As such, they are compressed in the same manner that integer types would be. However, since the values would be larger than they appear (because of the value 10 being stored as 10.000, or 10000 in the physical storage), the compression would be less than for an integer of the same magnitude.

Approximate Numeric Data

Approximate numeric values contain a decimal point and are stored in a format that's fast to manipulate. They are called *floating point* because they have a fixed number of significant digits, but the placement of the decimal point "floats," allowing for really small numbers or really large numbers. Approximate numeric values have some important advantages, as you'll see later in this appendix.

Approximate is such a negative term, but it's technically the proper term. It refers to the `real` and `float` datatypes, which are IEEE 754 standard single- and double-precision floating point values. The number is stored as a 32-bit or 64-bit value, with four parts:

- *Sign*: Determines whether this is a positive or negative value.
- *Exponent*: The exponent in base-2 of the mantissa.
- *Mantissa*: Stores the actual number that's multiplied by the exponent (also known as the *coefficient* or *significand*).
- *Bias*: Determines whether the exponent is positive or negative.

A complete description of how these datatypes are formed is beyond the scope of this book but may be obtained from the IEEE body at <http://www.ieee.org>.

- `float [(N)]`
 - *Domain*: $-1.79E + 308$ through $1.79E + 308$. The `float` datatype allows you to specify a certain number of bits to use in the mantissa, from 1 to 53. You specify this number of bits with the value in `N`. The default is 53.
 - *Storage*: See Table B-1.
- `real`
 - `real` is a synonym for `float(24)`.

Table B-1. *Floating Point Precision and Storage Requirements*

N (Number of Mantissa Bits for Float)		Precision	Storage Size
1–24	7	4 bytes	
25–53	15	8 bytes	

At this point, SQL Server rounds all values of `N` up to either 24 or 53. This is the reason that the storage and precision are the same for each of the values.

Note There's another approximate datatype named `real`. It's a synonym for a `float(24)` datatype. It can store values in the range from $-3.40E + 38$ through $3.40E + 38$.

Discussion:

Using these datatypes, you can represent most values from $-1.79E + 308$ to $1.79E + 308$ with a maximum of 15 significant digits. This isn't as many significant digits as the numeric datatypes can deal with, but the range is enormous and is plenty for almost any scientific application. These datatypes have a *much* larger range of values than any other datatype. This is because the decimal point isn't fixed in the representation. In numeric types, you always have a pattern such as `NNNNNN.NNNNN` for numbers. You can't store more digits than this to the left or right of the decimal point. However, with `float` values, you can have values that fit the following patterns (and much larger):

- `0.DDDDDDDDDDDDDDD`
- `NNNN.DDDDDDDDD`

- 0.000000000000000000000000000000DDDDDDDDDDDDDD
- NNNNNNNNNNNNNNNN000000000000000000

So, you have the ability to store tiny numbers, or large ones. This is important for scientific applications where you need to store and do math on an extreme range of values. The `float` datatypes are well suited for this usage.

Row Compression Effect:

The least significant bytes with all zeros are not stored. This is applicable mostly to non-fractional values in the mantissa.

Date and Time Data

There are two datatypes for working with date and time values: `datetime` and `smalldatetime`. Both have a time element and a date element, and you cannot separate them. Not having a simple date or time datatype can be a real bother at times, because often we want to store just the time of an event, or just the date of the event. For date-only values, it's simple: we just ignore the time by making sure that the time value is exactly midnight, and it works OK. Time values can be more of an issue. Although we can set the date to some arbitrary date and just use the time value, the date value then looks funny when viewed without formatting. We often build our own datatype for these values. I'll discuss this in this section as well.

date

Domain: Date-only values from January 1, 0001, to December 31, 9999.

Storage: 3-byte integer, storing the offset from January 1, 0001.

Accuracy: One day.

Discussion:

Of all the features added to SQL Server in 2008, this one datatype is worth the price of the upgrade (especially since I don't have to whip out my wallet and pay for it). The problem of how to store date values without time has plagued T-SQL programmers since the beginning of time (a.k.a. version 1.0).

With this type, you will be able to avoid the goofy tricks you have needed to go through to ensure that date types had no time in order to store just a date value. In the past versions of the book, I have even advocated the creation of your own pseudodate values stored in an integer value. That worked, but you could not use the built-in date functions without doing some “tricks.”

Row Compression Effect:

Technically you get the same compression as for any integer value, but dates in the “normal” range of dates require 3 bytes, meaning no compression is realized.

time [(precision)]

Domain: Time of day only (note this is not a quantity of time, but a point in time on the clock).

Storage: 3–5 bytes, depending on precision.

Accuracy: To 100 nanoseconds, depending on how it is declared. `time(0)` is accurate to 1 second, `time(1)` to .1 seconds, up to `time(7)` as .0000001. The default is 100-nanosecond accuracy.

Discussion:

Where the date type was earth-shattering, the time type is handy to have but generally less useful. It will seem like a good idea to store a point in time, but in that case you will have to make sure that both times are for the same day. Rather, for the most part when you want to store a time value, it is a point in time, and you need one of the date + time types.

The time value can be useful for storing a time for a recurring activity, for example, where the time is for multiple days rather than a single point in time.

Row Compression Effect:

Technically you get the same compression as for any integer value, but time values generally use most of the bytes of the integer storage, so very little compression should be expected for time values.

smalldatetime

Domain: Date and time data values between January 1, 1900, and June 6, 2079.

Storage: 4 bytes (two 2-byte integers: one for the day offset from January 1, 1900, the other for the number of minutes past midnight).

Accuracy: One minute.

Discussion:

The `smalldatetime` datatype is accurate to 1 minute. It requires 4 bytes of storage. `smalldatetime` values are the best choice when you need to store just the date, and possibly the time, of some event where accuracy of a minute isn't a problem.

Row Compression Effect:

When there is no time stored, 2 bytes can be saved, and times less than 4 a.m. can save 1 byte.

datetime

Domain: Date and time data values between January 1, 1753, and December 31, 9999.

Storage: 8 bytes (two 4-byte integers: one for the day offset from January 1, 1753, and the other for the number of 3.33-millisecond periods past midnight).

Accuracy: 3.33 milliseconds.

Discussion:

Using 8 bytes, `datetime` does require a sizable chunk of memory. There are few cases where you need this kind of precision. A primary example of where you do need this kind of precision is a timestamp column (not to be confused with the `timestamp` datatype, which I'll discuss in a later section), used to denote exactly when an operation takes place. This isn't uncommon if you need to get timing information, such as the time taken between two activities in seconds, or if you want to use the `datetime` value in a concurrency control mechanism.

Row Compression Effect:

For the date part of the type, dates before 2079 can save 1 byte. For the time part, 4 bytes are saved when there is no time saved, and it uses the first 2 bytes after the first 2 minutes and reaches the fourth byte after 4 a.m. After 4 a.m., compression can generally save 1 byte.

datetime2 [(precision)]

Domain: Dates from January 1, 0001, to December 31, 9999, with a time component.

Storage: Between 6 and 8 bytes. The first 4 bytes are used to store the date, and the others an offset from midnight, depending on the accuracy.

Accuracy: To 100 nanoseconds, depending on how it is declared. `datetime(0)` is accurate to 1 second, `datetime(1)` to .1 seconds, up to `datetime(7)` as .0000001. The default is 100-nanosecond accuracy.

Discussion:

`datetime2`: Despite the hideous Oracle-looking name, this is a much better datatype than `datetime`. Technically you get far better time support without being limited by the .003 accuracy issues that `datetime` is. The only downside I currently see is support for the type in your API.

What I see as the immediate benefit of this type is to fix the amount of accuracy that your users actually desire. Most of the time a user doesn't desire fractional seconds, unless the purpose of the type is something scientific or technical. With `datetime2`, you can choose 1-second accuracy. Also, you can store .999 seconds, unlike `datetime`, which would round .999 up to 1, whereas .998 would round down to .997.

Row Compression Effect:

For the date part of the type, dates before 2079 can save 1 byte of the 4 bytes for the date. Little compression should be expected for the time portion.

datetimeoffset [(precision)]

The `datetimeoffset` is the same as `datetime2`, but it includes an offset from UTC time (8–10 bytes).

Domain: Dates from January 1, 0001, to December 31, 9999, with a time component. Includes the offset from the UTC time, in a format of `[+|-] hh:mm`. (Note that this is not time zone/day-light saving time aware. It simply stores the offset at the time of storage.)

Storage: Between 8 and 10 bytes. The first 4 bytes are used to store the date, and just like `datetime2`, 2–4 will be used for the time, depending on the accuracy. The UTC offset is stored in the additional 2 bytes.

Accuracy: To 100 nanoseconds, depending on how it is declared. `datetimeoffset(0)` is accurate to 1 second, `datetimeoffset (1)` to .1 seconds, up to `datetimeoffset (7)` as .0000001. The default is 100-nanosecond accuracy.

Discussion:

The offset seems quite useful but in many ways is more cumbersome than using two date columns, one for UTC and one for local (though this will save a bit of space). Its true value is that it defines an exact point in time better than the regular `datetime` type, since there is no ambiguity as to where the time was stored.

A useful operation is to translate the date from its local offset to UTC, like this:

```
DECLARE @LocalTime DateTimeOffset
SET @LocalTime = SYSDATETIMEOFFSET()
SELECT @LocalTime
SELECT SWITCHOFFSET(@LocalTime, '+00:00') As UTCTime
```

The true downside is that it stores an offset, not the time zone, so daylight saving time will still need to be handled manually.

Row Compression Effect:

For the date part of the type, dates before 2079 can save 1 byte of the 4 bytes for the date. Little compression should be expected for the time portion.

Discussion on All Date Types

Date types are often some of the most troublesome types for people to deal with. In this section, I'll lightly address the following problem areas:

- Date functions
- Date ranges
- Representing dates in text formats

Date Functions

With the creation of new date and time (and datetime) types in SQL Server 2008, there needed to be more functions to work with. Microsoft has added functions that return and modify dates and time with more precision than GETDATE or GETUTCDATE:

- SYSDATETIME: Returns system time to the nearest fraction of a second, with seven places of scale in the return value
- SYSDATETIMEOFFSET: Same as SYSDATETIME but returns the offset of the server in the return value
- SYSUTCDATETIME: Same as SYSDATETIME but returns the UTC date time rather than local time
- SWITCHOFFSET: Changes the offset for a datetimeoffset value
- TODATETIMEOFFSET: Converts a local date and time value to a given time zone

There have been a few changes to the basic date functions as well:

- DATENAME: Includes values for microsecond, nanosecond, and TZoffset. These will work differently for the different types. For example, TZoffset will work with datetime2 and datetimeoffset, but not the other types.
- DATEPART: Includes microsecond, nanosecond, TZoffset, and ISO_WEEK. ISO_WEEK is a feature that has been long desired by programmers who need to know the week of the year, rather than the nonstandard week value that SQL Server has previously provided.
- DATEADD: Supports micro and nanoseconds.

With all the date functions, you really have to be careful that you are cognizant of what you are requesting. For an example, how old is this person? Say you know that a person was born on December 31, 2008, and on January 3, 2009, you want to know their age (yes, this seems like a very easy question, but the answer is kind of interesting). Common sense says to look for a function to take the difference between two dates. You are in luck—there is a DATEDIFF function. Executing the following:

```
DECLARE @time1 date = '20081231',
        @time2 date = '20090102'
SELECT DATEDIFF(yy,@time1,@time2)
```

You see that the person is 0 years old right? Wrong! It returns 1. Well, maybe. But if that is true, then the following should probably return 2, right?

```
DECLARE @time1 date = '20080101',
        @time2 date = '20091231'
SELECT DATEDIFF(yy,@time1,@time2)
```

No, shucks, that also returns 1. So, no matter whether the date values are 1 day apart or 730, you get the same result? Yes, because the DATEDIFF function is fairly dumb in that it is taking the difference of the year value, not the difference in years. Then, to find out the age, you will need to use several functions.

So, what is the answer? We could do something more fancy, likely by coming up with an algorithm based on the number of days, or months, or shifting dates to some common value, but that is way more like work than the really straightforward answer. Build a table of dates, commonly called a *calendar* (a calendar table is introduced in Chapter 7).

Date Ranges

This topic will probably seem really elementary, but the fact is that one of the largest blunders in the database implementation world is working with ranges of dates. The problem is that when you want to do inclusive ranges, you have always needed to consider the time in the equation. For example, the following criteria:

```
WHERE pointInTimeValue between '2008-01-01' and '2008-12-31'
```

means something different based on whether the values stored in dateValue have, or do not have, a time part stored. With the introduction of the date type, the need to worry about this issue of date ranges will probably become a thing of the past, but it is still an issue that you'll need to worry about in the near-term at least.

Thinking back to the example WHERE clause, the problem is that any value with the same date as the end value plus a time (such as '2008-12-31 12:00:00') does not fit within the preceding selecting criteria. So, this value will be left out of the results. Of course, the next criteria to try would logically be the following:

```
WHERE pointInTimeValue between '2009-01-01' and '2009-12-31'
```

But now you have actually missed all the activity that occurred on December 31 that wasn't at midnight. And of course, this problem gets worse when you move toward smaller ranges of time.

There are two ways to deal with this. Either code your WHERE clause like this:

```
WHERE pointInTimeValue >= '2009-01-01' and pointInTimeValue < '2010-01-01'
```

or use a calculated column to translate point-in-time values in your tables to date-only values (like dateValue as cast(pointInTimeValue as date)). Having done that, a value such as '2008-12-31 12:00:00' will be truncated to '2008-12-31 00:00:00', and a row containing that value will be picked up by selection criteria such as this:

```
WHERE pointInTimeValue between '2008-01-01' and '2008-12-31'
```

A common solution that I don't generally suggest is to use a between range like this:

```
WHERE pointInTimeValue between '2009-01-01' and '2009-12-31 23:59:59.9999999'
```

The idea is that if the second value is less than the next day, then values for the next day won't be returned. The major problem with this solution has to do with the conversion of 23:59:59.9999999 to some date datatype. For example, if `pointInTimeValue` were a `Datetime`, it would be rounded up to the next day because the maximum time for a `datetime` is 23:59:59.997. For the range to work with a `datetime`, you would use this:

```
WHERE pointInTimeValue between '2009-01-01' and '2009-12-31 23:59:59.997'
```

However, for a `smalldatetime` value, it would need to be this:

```
WHERE pointInTimeValue between '2009-01-01' and '2009-12-31 23:59'
```

and so on, for all of the different date types, which gets complicated by the new types where you can specify precision. We strongly suggest you avoid trying to use a maximum date value like this unless you are tremendously careful with the types of data and how their values round off.

For more information about how date and time data work with one another and converting from one type to another, read the topic "Using Date and Time Data" in Books Online.

Representing Dates in Text Formats

When working with date values in text, using a standard format is always best. There are many different formats used around the world for dates, most confusingly `MMDDYYYY` and `DDMMYYYY` (is 01022004 or 02012004 the same day, or a different day?). Although SQL Server uses the locale information on your server to decide how to interpret your date input, using one of the following formats ensures that SQL Server doesn't mistake the input regardless of where the value is entered.

Generally speaking, it is best to stick with one of the standard date formats that are recognized regardless of where the user is. This prevents any issues when sharing data with international clients, or even with sharing it with others on the Web when looking for help.

There are several standards formats that will work:

- ANSI SQL Standard
 - *No time zone offset*: 'YYYY-MM-DD HH:MM:SS'
 - *With time zone*: 'YYYY-MM-DD HH:MM:SS -OH:OM' (Z can be used to indicate the time zone is 00:00)
- ISO 8601
 - *Unseparated*: 'YYYYMMDD'
 - *Numeric*: 'YYYY-MM-DD'
 - *Time*: 'HH:MM:SS.ssssss' (SS and .sssssss are optional)
 - *Date and time*: 'YYYY-MM-DDTHH:MM:SS.ssssss'
 - *Date and time with offset*: 'YYYY-MM-DDTHH:MM:SS.ssssss -OH:OM'
- ODBC
 - *Date*: {d 'YYYY-MM-DD'}
 - *Time*: {t 'HH:MM:SS'}
 - *Date and time*: {ts 'YYYY-MM-DD HH:MM:SS'}

Using the ANSI SQL Standard or the ISO 8601 formats is generally considered the best practice for specifying date values. It will definitely feel odd when you first begin typing '2008-08-09' for a date value, but once you get used to it, it will feel natural.

The following are some examples using the ANSI and ISO formats:

```
select cast ('2009-01-01' as smalldatetime) as dateOnly
select cast('2009-01-01 14:23:00.003' as datetime) as withTime
```

You might also see values that are close to this format, such as the following:

```
select cast ('20090101' as smalldatetime) as dateOnly
select cast('2009-01-01T14:23:00.120' as datetime) as withTime
```

These are acceptable variations. For more information, check SQL Server 2008 Books Online under “Using Date and Time Data.”

Character Strings

Most data that's stored in SQL Server uses character datatypes. In fact, usually far too much data is stored in character datatypes. Frequently, character columns are used to hold noncharacter data, such as numbers and dates. Although this might not be technically wrong, it isn't ideal. For starters, storing a number with eight digits in a character string requires at least 8 bytes, but as an integer it requires 4 bytes. Searching on integers is far easier because 1 always precedes 2, whereas 11 comes before 2 in character strings. Additionally, integers are stored in a format that can be manipulated using intrinsic processor functions, as opposed to having SQL Server-specific functions deal with the data.

char[(length)]

Domain: ASCII characters, up to 8,000 characters long.

Storage: 1 byte * length.

Discussion:

The char datatype is used for fixed-length character data. Every value will be stored with the same number of characters, up to a maximum of 8,000 bytes. Storage is exactly the number of bytes as per the column definition, regardless of actual data stored; any remaining space to the right of the last character of the data is padded with spaces. The default size if not specified is 1 (it is best practice to include the size).

You can see the possible characters by executing the following query:

```
SELECT number, CHAR(number)
FROM   utility.sequence
WHERE  number >=0 and number <= 255
```

Tip The sequence table is a common table that every database should have. It's a table of integers that can be used for many utility purposes. In Chapter 7, I present a sequence table that you can use for this query.

The maximum limit for a char is 8,000 bytes, but if you ever get within a mile of this limit for a fixed-width character value, you're likely making a big design mistake because it's very rare to have massive character strings of exactly the same length. You should employ the char datatype only in cases where you're guaranteed to have exactly the same number of characters in every row.

The char datatype is most often used for codes and identifiers, such as customer numbers or invoice numbers where the number includes alpha characters as well as integer data. An example is a vehicle identification number (VIN), which is stamped on most every vehicle produced around the world. Note that this is a composite attribute, because you can determine many things about the automobile from its VIN.

Another example where a char column is usually found is in Social Security numbers (SSNs), which always have nine characters and two dashes embedded.

Row Compression Effect:

Instead of storing the padding characters, it removes them for storage and adds them back whenever the data is actually used.

Note The setting ANSI_PADDING determines exactly how padding is handled. If this setting is ON, the table is as I've described; if not, data will be stored as I'll discuss in the "varchar[(length)]" section that follows. It's best practice to leave this ANSI setting ON.

varchar[(length)]

Domain: ASCII characters, up to 8,000 characters long.

Storage: 1 byte * length + 2 bytes (for overhead).

Discussion:

For the varchar datatype, you choose the maximum length of the data you want to store, up to 8,000 bytes. The varchar datatype is far more useful than char, because the data doesn't have to be of the same length and SQL Server doesn't pad out excess memory with spaces. There's some reasonably minor overhead in storing variable-length data. First, it costs an additional 2 bytes per column. Second, it's a bit more difficult to get to the data, because it isn't always in the same location of the physical record. The default size if not specified is 1 (it is best practice to include the size).

Use the varchar datatype when your character data varies in length. The good thing about varchar columns is that, no matter how long you make the maximum, the space used by the column is based on the actual size of the characters being stored plus the few extra bytes that specify how long the data is.

You'll generally want to choose a maximum limit for your datatype that's a reasonable value, large enough to handle most situations, but not too large as to be impractical to deal with in your applications and reports. For example, take people's first names. These obviously require the varchar type, but how long should you allow the data to be? First names tend to be a maximum of 15 characters long, though you might want to specify 20 or 30 characters for the unlikely exception.

The most prevalent storage type for non-key values that you'll use is varchar data, because, generally speaking, the size of the data is one of the most important factors in performance tuning. The smaller the amount of data, the less has to be read and written. This means less disk access, which is one of the two most important bottlenecks we have to deal with (networking speed is the other).

Row Compression Effect:

No effect.

varchar(max)

Domain: ASCII characters, up to $2^{31} - 1$ characters (that is a maximum of 2GB worth of text!).

Storage: There are a couple possibilities for storage based on the setting of the table option 'large value types out of row', which is set with the `sp_tableoption` system stored procedure:

- OFF or 0 =: The data for all the columns fits in a single row, and the data is stored in the row with the same storage costs for non-max varchar values. Once the data is too big to fit in a single row, data can be placed on more than one row. This is the default setting.
- ON or 1 =: You store varchar(max) values using 16-byte pointers to separate pages outside the table. Use this setting if the varchar(max) data will only seldom be used in queries.

Discussion:

The varchar(max) datatype is possibly the greatest thing since someone said, “Meat and fire—I wonder if they might go together?” Too long we struggled with the painful text datatype and all its quirks. You can deal with varchar(max) values using mostly the same functions and methods that you use with normal varchar values. There’s a minor difference, though. As the size of your varchar(max) column grows toward the upper boundary, it’s likely true that you aren’t going to want to be sending the entire value back and forth over the network most of the time. I know that even on my 100MB LAN, sending 2GB is no instantaneous operation, for sure.

There are a couple things to look at, which I’ll just touch on here:

- The UPDATE statement has a .WRITE() clause to write chunks of data to the (max) datatypes. This is also true of varbinary(max).
- Unlike text and image values, (max) datatypes are accessible in AFTER triggers.

One word of warning for when your code mixes normal varchar and varchar(max) values in the same statement: normal varchar values do not automatically change datatype to a (max) type when the data being manipulated grows beyond 8,000 characters. For example, write a statement such as the following:

```
DECLARE @value varchar(max)
SET @value = replicate('X',8000) + replicate('X',8000)
SELECT len(@value)
```

This returns the following result, which you would expect to be 16000, since you have two 8,000-character strings:

```
-----
8000
```

The reason is that the type of the REPLICATE function is varchar, when replicating normal char values. Adding two varchar values together doesn’t result in a varchar(max) value. However, most of the functions return varchar(max) values when working with varchar(max) values. For example:

```
DECLARE @value varchar(max)
SET @value = replicate(cast('X' as varchar(max)),16000)
SELECT len(@value)
```

This returns the following result:

16000

Row Compression Effect:

No effect.

text

Don't use the text datatype for any reason in new designs. It might not exist in the next version of SQL Server. Replace with `varchar(max)` whenever you possibly can. See SQL Server Books Online for more information.

Unicode Character Strings: `nchar`, `nvarchar`, `nvarchar(max)`, `ntext`

Domain: ASCII characters, up to $2^{15} - 1$ characters (2GB of storage).

Storage: Same as other character datatypes, though every character takes 2 bytes rather than 1. (Note there is no support for any of the variable-length Unicode storage.)

Discussion:

So far, the character datatypes we've been discussing have been for storing typical ASCII data. In SQL Server 7.0 (and NT 4.0), Microsoft implemented a new standard character format called Unicode. This specifies a 16-bit character format that can store characters beyond just the Latin character set. In ASCII—a 7-bit character system (with the 8 bits for Latin extensions)—you were limited to 256 distinct characters. This was fine for most English-speaking people but was insufficient for other languages. Asian languages have a character for each different syllable and are nonalphabetic; Middle Eastern languages use several different symbols for the same letter according to its position in the word. Unicode expanded the amount of characters and eliminated the need for code pages to allow for a vastly expanded character set (which allowed you to have multiple character sets in an 8-character encoding set in ASCII). SQL Server supports the Unicode Standard, version 3.2.

For these datatypes, you have the `nchar`, `nvarchar`, `nvarchar(max)`, and `ntext` datatypes. They are the same as the similarly named types (without the `n`) that we've already described, except for one thing: Unicode uses double the number of bytes to store the information, so it takes twice the space, thus cutting by half the number of characters that can be stored.

One quick tip: if you want to specify a Unicode value in a string, you append an `N` to the front of the string, like so:

```
SELECT N'Unicode Value'
```

Tip You should migrate away from `ntext` as a datatype just as you should for the `text` datatype.

Row Compression Effect:

Same as their ASCII counterparts. This has no effect on the variable types and doesn't store trailing blanks for the fixed-length types.

Binary Data

Binary data allows you to store a string of bytes. It's useful for storing just about anything, especially data from a client that might or might not fit into a character or numeric datatype. In SQL Server 2005, binary columns have become even more useful, because you can use them when storing encrypted data. In Chapter 7, you'll learn about the encryption capabilities of SQL Server 2005.

One of the restrictions of binary datatypes is that they don't support bitwise operators, which would allow you to do some powerful bitmask storage by being able to compare two binary columns to see not only whether they differ, but how they differ. The whole idea of the binary datatypes is that they store strings of bits. The bitwise operators can operate on integers, which are physically stored as bits. The reason for this inconsistency is fairly clear from the point of view of the internal query processor. The bitwise operations are operations that are handled in the processor, whereas the binary datatypes are SQL Server specific.

Binary literal values are specified as 0xB1B2B3 . . . BN. 0x tells you that it's a hexadecimal value. B1 specifies the first single byte in hexadecimal.

binary[(length)]

Domain: Fixed-length binary data with a maximum length of 8,000 bytes.

Storage: Number of bytes the value is defined for. The default length is 1, if not specified (it is best practice to include a size).

Discussion:

The use of binary columns is fairly limited. You can use them to store any binary values that aren't dealt with by SQL Server. Data stored in binary is simply a string of bytes:

```
declare @value binary(10)
set @value = cast('helloworld' as binary(10))
select @value
```

This returns the following result:

```
-----
0x68656C6C6F776F726C64
```

Now you can reverse the process:

```
select cast(0x68656C6C6F776F726C64 as varchar(10))
```

This returns the following result:

```
-----
helloworld
```

Note that casting the value HELLOWORLD gives you a different value:

```
-----
0x48454C4C4F574F524C44
```

This fact that these two binary values are different, even for textual data that would be considered equivalent on a case-insensitive collation, has been one use for the binary datatype: case-sensitive searches. It's far more efficient to use the COLLATE keyword and use a different collation if you want to do a case-insensitive comparison on text data.

Row Compression Effect:

Trailing zeros are not stored but are returned when the values are used.

varbinary[(length)]

Domain: Variable-length binary data with a maximum length of 8,000 bytes.

Storage: Number of bytes the value is defined for, plus 2 bytes for variable-length overhead. The default length is 1, if not specified (it is a best practice to include a size).

Discussion:

The usage is the same as binary, except the number of bytes is variable.

Row Compression Effect:

No effect.

varbinary(max)

Domain: Binary data, up to $2^{31} - 1$ bytes (up to 2GB for storage) when data is stored in SQL Server files, up to the max of the storage for data stored in the filestream. For more information and examples about the filestream, check Chapter 7.

Storage: There are a couple possibilities for storage based on whether the data is stored using the filestream setting, as well as the setting of the table option 'large value types out of row':

- OFF =: If the data for all the columns fits in a single row, the data is stored in the row with the same storage costs for non-max varchar values. Once the data is too big to fit in a single row, data can be placed on greater than one row.
- ON =: You store varbinary(max) values using 16-byte pointers to separate pages outside the table. Use this setting if the varchar(max) data will only seldom be used in queries.

Discussion:

The varbinary(max) datatype provides the same kinds of benefits for large binary values as the varchar(max) does for text. Pretty much you can deal with varbinary(max) values using the same functions and the same methods as you do with the normal varbinary values.

What's cool is that you can store text, JPEG and GIF images, and even Word documents and Excel spreadsheet data using the varbinary(max) type. On the other hand, it can be much slower and more programming work to use SQL Server as a storage mechanism for files, mostly because it's slow to retrieve really large values from the database as compared to from the file system. You can, however, use a filestream access to get the best of both possible worlds by using Win32 access to a file in a directory within the context of a transaction. This approach is described in greater detail in Chapter 7.

Row Compression Effect:

No effect.

image

Just like the text datatype, the image datatype has always been a real bother and is being deprecated in this version of SQL Server. Don't use the image datatype in new designs if at all possible. It very well may not exist in the next version of SQL Server. Replace with `varbinary(max)` in any location you can. See SQL Server Books Online for more information or if you have existing image column data that you need to manipulate.

Other Datatypes

The following datatypes are somewhat less easy to categorize but are still commonly employed in OLTP systems:

- `rowversion (timestamp)`
- `uniqueidentifier`
- `cursor`
- `table`
- `sql_variant`

rowversion (a.k.a. timestamp)

The rowversion datatype is a database-wide unique number. When you have a rowversion column in a table, the value of the rowversion column changes for each modification to each row. The value in the rowversion column is guaranteed to be unique across all tables in the datatype. It's also known as a timestamp value, but it doesn't have any time implications—it's merely a unique value to tell you that your row has changed.

Tip In the SQL standards, a timestamp datatype is equivalent to what you know as a datetime datatype. To avoid confusion, Microsoft now recommends that you specify rowversion rather than timestamp.

The rowversion column of a table (you may have only one) is usually used as the data for an optimistic locking mechanism. The rowversion datatype is a mixed blessing. It's stored as an 8-byte varbinary value. Binary values aren't always easy to deal with, and their use depends on which mechanism you're using to access your data.

As an example of how the rowversion datatype works, consider the following batch:

```
SET nocount on
CREATE TABLE testRowversion
(
    value    varchar(20) NOT NULL,
    auto_rv  rowversion NOT NULL
)

INSERT INTO testRowversion (value) values ('Insert')

SELECT value, auto_rv FROM testRowversion
UPDATE testRowversion
SET value = 'First Update'
```

```

SELECT value, auto_rv from testRowversion
UPDATE testRowversion
SET value = 'Last Update'

```

```

SELECT value, auto_rv FROM testRowversion

```

This batch returns the following results:

value	auto_rv
Insert	0x000000000000007D1
First Update	0x000000000000007D2
Last Update	0x000000000000007D3

You didn't touch the `auto_rv` column, and yet it incremented itself twice. However, you can't bank on the order of the `rowversion` values being sequential, because updates of other tables will change the value as well. All `rowversion` values in a database draw from the same pool of values. It's also in your best interest not to assume in your code that a `rowversion` number is an incrementing value. How `rowversions` are implemented is a detail that will likely change in the future. If a better method of building database-wide unique values comes along that's even a hair faster, Microsoft will likely use it.

You can create variables of the `rowversion` type for holding `rowversion` values, and you can retrieve the last-used `rowversion` via the `@@dbts` configuration function. `Rowversion` columns are used in Chapter 9, when I demonstrate optimistic locking.

Row Compression Effect:

Uses an integer representation of the value, using 8 bytes. Then it can be compressed just like the `bigint` type.

uniqueidentifier

Globally unique identifiers are fast becoming a mainstay of Microsoft computing. The name says it all—these identifiers are globally unique. According to the way that GUIDs are formed, there's a tremendously remote chance that there will ever be any duplication in their values. They're generated by a formula that includes the current date and time, a unique number from the CPU clock, and some other “magic numbers.”

In your databases, these GUID values are stored in the `uniqueidentifier` type. An interesting use is to have a key value that's guaranteed to be unique across databases and servers. You can generate a GUID value in T-SQL using the `newid` function.

```

DECLARE @guidVar uniqueidentifier
SET @guidVar = newid()

SELECT @guidVar as guidVar

```

returns

guidVar

6C7119D5-D48F-475C-8B60-50D0C41B6EBF

GUIDs are stored as 16-byte binary values. Note that a GUID isn't exactly a straight 16-byte binary value. You cannot put just any binary value into a `uniqueidentifier` column, because the value must meet the criteria for the generation of a GUID, which aren't well documented, for obvious reasons. (For more information, a good resource is <http://en.wikipedia.org/wiki/guid>.)

If you need to create a `uniqueidentifier` column that's autogenerating, you can set a property in the `CREATE TABLE` statement (or `ALTER TABLE`, for that matter). It's the `rowguidcol` property, and it's used like so:

```
CREATE TABLE guidPrimaryKey
(
    guidPrimaryKeyId uniqueidentifier NOT NULL rowguidcol DEFAULT newId(),
    value varchar(10)
)
```

I've introduced a couple new things here: `rowguidcol` and default values. Suffice it to say that if you don't provide a value for a column in an insert operation, the default operation will provide it. In this case, you use the `newId()` function to get a new `uniqueidentifier`. Execute the following `INSERT` statement:

```
INSERT INTO guidPrimaryKey(value)
VALUES ('Test')
```

Then run the following command to view the data entered:

```
SELECT *
FROM guidPrimaryKey
```

This returns the following result (though of course your key value will be different):

guidPrimaryKeyId	value
8A57C8CD-7407-47C5-AC2F-E6A884C7B646	Test

The `rowguidcol` property of a column built with the `uniqueidentifier` notifies the system that this is just like an identity column value for the table—a unique pointer to a row in a table. Note that neither the identity nor the `rowguidcol` properties guarantee uniqueness. To provide such a guarantee, you have to implement your tables using `UNIQUE` constraints.

It would seem that the `uniqueidentifier` would be a better way of implementing primary keys, because when they're created, they're unique across all databases, servers, and platforms. However, there are two main reasons why you won't use `uniqueidentifier` columns to implement all your primary keys:

- *Storage requirements:* Because they're 16 bytes in size, they're considerably more bloated than a typical integer column.
- *Typeability:* Because there are 36 characters in the textual version of the GUID, it's hard to type the value of the GUID into a query, and it isn't easy to enter.

If you're using the GUID values for the primary key of a table and you're clustering on this value, you can use another function to generate the values: `newSequentialId()`. You can use this

function only in a default constraint. It's used to guarantee that the next GUID chosen will be greater than the previous value:

```

DROP TABLE guidPrimaryKey
go
CREATE TABLE guidPrimaryKey
(
    guidPrimaryKeyId uniqueidentifier NOT NULL
                        rowguidcol DEFAULT newSequentialId(),
    value varchar(10)
)
GO
INSERT INTO guidPrimaryKey(value)
SELECT 'Test'
UNION ALL
SELECT 'Test1'
UNION ALL
SELECT 'Test2'
GO

SELECT *
FROM guidPrimaryKey

```

This returns something like the following:

guidPrimaryKeyId	value
-----	-----
AA52457C-339B-DA11-9A3C-001422E6CCC3	Test
AB52457C-339B-DA11-9A3C-001422E6CCC3	Test1
AC52457C-339B-DA11-9A3C-001422E6CCC3	Test2

Note You may notice that the increasing value appears to be in the letters to the far left. To the naked eye, it would appear that we could be pretty close to running out of values, since the progression of AA, AB, AC would run out pretty quickly. The fact is, the values are not being sorted on the text representation of the GUID, but on the internal binary value. If you are interested in learning more about the internals of GUIDS, check <http://en.wikipedia.org/wiki/guid>.

Now, using a GUID for a primary key is just about as good as using an identity column for building a surrogate key, particularly one with a clustered index (they are still rather large at 16 bytes versus 4 for an integer, or even 8 for a bigint). That's because all new values will be added to the end of the index rather than randomly throughout the index. (Chapter 8 covers indexes, but be cognizant that a random value distributed throughout your rows can cause fragmentation unless you provide a fill factor that allows for adding rows to pages.) Values in the `uniqueidentifier` type will still be four times as large as an integer column, hence requiring four times the storage space. This makes using a `uniqueidentifier` a less than favorable index candidate from the database storage layer's perspective. However, the fact that it can be generated by any client and be guaranteed unique is a major plus, rather than requiring you to generate them in a single threaded manner to ensure uniqueness.

Row Compression Effect:

No effect.

cursor

A cursor is a mechanism that allows row-wise operations instead of using the normal set-wise way. You use the `cursor` datatype to hold a reference to a SQL Server T-SQL cursor. You may not use a `cursor` datatype as a column in a table. Its only use is in T-SQL code to hold a reference to a cursor, which can be passed as a parameter to a stored procedure.

Row Compression Effect:

No effect.

table

The table type is kind of two different things now in 2008. First you have the table type that is essentially a temporary table that you can declare like a variable at runtime, and you can define its characteristics. Second (and new to 2008), you have table types that are defined and stored for later use, for example, as table-valued parameters. I have broken these two different types of uses down into two sections. Neither usage is affected by row compression.

Table Variables

The table variable has a few things in common with the `cursor` datatype, but instead of a cursor, it holds a reference to a result set. The name of the datatype is a pretty bad choice, because it will make functional programmers think that they can store a pointer to a table. It's actually used to store a result set as a temporary table. In fact, the table is exactly like a temporary table in implementation. However, you don't get any kind of statistics on the table, nor are you able to index the table datatype, other than to apply `PRIMARY KEY` and `UNIQUE` constraints in the table declaration. You can also have `CHECK` and `DEFAULT` constraints.

Unlike local temporary tables (those declared with `#` preceding the name), table datatype variables won't cause recompiles in stored procedures that use them, because they don't have any statistics to change the plan anyway. Use them only for modestly small sets of data (hundreds of rows, not thousands, generally), such as when all the data in the table can fit on a single data page.

The following is an example of the syntax needed to employ the table variable type:

```
DECLARE @tableVar TABLE
(
    id int IDENTITY PRIMARY KEY,
    value varchar(100)
)
INSERT INTO @tableVar (value)
VALUES ('This is a cool test')

SELECT id, value
FROM @tableVar
```

This returns the following result:

id	value
1	This is a cool test

As with the `cursor` datatype, you may not use the table datatype as a column in a table, and it can be used only in T-SQL code to hold a set of data. One of the primary purposes for the table datatype is for returning a table from a user-defined function, as in the following example:

```
CREATE FUNCTION table$testFunction
(
    @returnValue varchar(100)
)
RETURNS @tableVar table
(
    value varchar(100)
)
AS
BEGIN
    INSERT INTO @tableVar (value)
    VALUES (@returnValue)

    RETURN
END
```

Once created, you can use the table datatype returned by the function using typical SELECT syntax:

```
SELECT *
FROM dbo.table$testFunction('testValue')
```

This returns the following result:

value
testValue

One interesting thing about the table datatype is that it isn't subject to transactions. For example:

```
DECLARE @tableVar TABLE
(
    id int IDENTITY,
    value varchar(100)
)
BEGIN TRANSACTION
    INSERT INTO @tableVar (value)
    VALUES ('This will still be there')
ROLLBACK TRANSACTION

SELECT id, value
FROM @tableVar
```

This returns the following result:

id	value
1	This will still be there

For this reason, these tables are useful for logging errors, because the data is still available after the ROLLBACK TRANSACTION.

Table Valued Parameters

One of the oft-requested features for SQL Server was the ability to pass in a table of values to a stored procedure. Using the table type, you can now do this, but not in as free a manner as you

probably would have initially hoped. Instead of being able to define your table on the fly, you are required to use a type that you predefine.

The table type you will define is the same as the datatype alias we discussed in Chapter 5, except you specify an entire table, with all of the same things that a table variable can have, including PRIMARY KEY, UNIQUE, CHECK, and DEFAULT constraints.

An example that I imagine will be very commonly imitated is the generic table type with a list of integer values to pass as a parameter or to use in a query instead of an IN clause:

```
CREATE TYPE GenericIdList AS TABLE
(
    Id Int Primary Key
)
```

You declare the table variable just like any other and then load and use the variable with data just like any other local variable table:

```
DECLARE @ProductIdList GenericIdList

INSERT INTO @productIDList
VALUES (1),(2),(3),(4)

SELECT ProductID, Name, ProductNumber
FROM AdventureWorks2008.Production.product
    JOIN @productIDList as list
        on Product.ProductID = List.Id
```

This returns the following:

ProductID	Name	ProductNumber
1	Adjustable Race	AR-5381
2	Bearing Ball	BA-8327
3	BB Ball Bearing	BE-2349
4	Headset Ball Bearings	BE-2908

Of course, you can then use the type in your stored procedure creation statements as well:

```
CREATE PROCEDURE product$list
(
    @productIDList GenericIdList READONLY
)
AS
SELECT ProductID, Name, ProductNumber
FROM AdventureWorks2008.Production.product
    JOIN @productIDList as list
        on Product.ProductID = List.Id
```

Unfortunately, you cannot pass a set of row constructors to the stored procedure; instead, you will need to declare and load a table variable to use this construct from T-SQL.

```
DECLARE @ProductIdList GenericIdList

INSERT INTO @productIDList
VALUES (1),(2),(3),(4)

EXEC product$list @ProductIdList
```

What makes this really nice is that in ADO.NET, you can declare a `DataTable` object and pass it to the procedure as a parameter, just like any other value now. This will make the ability to insert multiple items at a time or `SELECT` multiple rows far easier than ever before. In the past, we used a kludgy comma-delimited list or XML to do this, and it worked, but not in a natural manner we are accustomed to, and it was generally slow. This method will now work in a natural manner, allowing us to finally support multiple operations in a single transaction from an easy-to-build ADO.NET construct.

sql_variant

The catchall datatype, the `sql_variant` type, allows you to store a value of almost any datatype that I've discussed. This ability allows you to create a column or variable where you don't know ahead of time exactly what kind of data will be stored. The `sql_variant` datatype allows you to store values of various SQL Server–supported datatypes, except for `varchar(max)`, `varbinary(max)`, `xml`, `text`, `ntext`, `rowversion/timestamp`, and `sql_variant`.

Note Although the `rowversion` datatype cannot be stored directly in a `sql_variant`, a `rowversion` value can be stored in a `binary(8)` variable, which can in turn be stored in a `sql_variant` variable. Also, it might seem strange that you can't store a variant in a variant, but this is just saying that the `sql_variant` datatype doesn't exist as such—SQL Server chooses the best type of storage in which to store the value you give to it.

Generally, `sql_variant` is a datatype to steer clear of unless you really cannot know the datatype of a given value until the user enters the value. I used the `sql_variant` in Chapter 7 when I implemented the user-specified data storage using the entity-attribute-value solution. This allowed the user to enter any type of data and then have the system store the data in the most appropriate method.

The `sql_variant` type has some obvious value, and I used it earlier in Chapter 7 when building an entity-attribute-value solution for an open schema solution. By not needing to know the type at design time, you can allow the user to insert any type of data that they might want.

However, the positives lead directly to the negatives to the `sql_variant` type. Although simple storage and viewing of the data isn't too hard, it isn't easy to manipulate data once it has been stored in a `sql_variant` column. I'll leave it to you to read the information fully in the parts of SQL Server Books Online that deal with variant data, but some issues to consider are as follows:

- *Difficulties assigning data from a `sql_variant` column to a stronger typed datatype:* You have to be careful, because the rules for casting a variable from one datatype to another are difficult and might cause errors if the data can't be cast. For example, you can't cast the `varchar(10)` value 'Not a Date' to a `datetime` datatype. Such problems become an issue when you start to retrieve the variant data out of the `sql_variant` datatype and try to manipulate it.
- *NULL `sql_variant` values are considered to have no datatype:* Hence, you'll have to deal with `sql_variant` NULLs differently from nulls in other datatypes.
- *Comparisons of variants to other datatypes could cause difficult-to-catch programmatic errors, because of the `sql_variant` value instance's datatype:* Usually, the compiler will know whether you try to run a statement that compares two incompatible datatypes, such as `@intVar = @varcharVar`. However, if the two variables in question were defined as `sql_variants` and the datatypes don't match, then the values won't match because of the datatype incompatibilities.

When working with `sql_variant` variables or columns, you can use the `SQL_VARIANT_PROPERTY` function to discover the datatype of a given `sql_variant` value. For example:

```
DECLARE @varcharVariant sql_variant
SET @varcharVariant = '1234567890'
SELECT @varcharVariant AS varcharVariant,
       SQL_VARIANT_PROPERTY(@varcharVariant, 'BaseType') as baseType,
       SQL_VARIANT_PROPERTY(@varcharVariant, 'MaxLength') as maxLength,
       SQL_VARIANT_PROPERTY(@varcharVariant, 'Collation') as collation
```

The preceding statement returns the following result:

VarcharVariant	baseType	maxLength	collation
1234567890	varchar	10	SQL_Latin1_General_CP1_CI_AS

For numeric data, you can also find the precision and scale:

```
DECLARE @numericVariant sql_variant
SET @numericVariant = 123456.789
SELECT @numericVariant AS numericVariant,
       SQL_VARIANT_PROPERTY(@numericVariant, 'BaseType') as baseType,
       SQL_VARIANT_PROPERTY(@numericVariant, 'Precision') as precision,
       SQL_VARIANT_PROPERTY(@numericVariant, 'Scale') as scale
```

This returns the following result:

numericVariant	baseType	precision	scale
123456.789	numeric	9	3

Not Simply Scalar Datatypes

This section will deal with the class of datatypes that have been implemented by Microsoft that aren't really scalar values. Another common term for these datatypes that have cropped up around the Internet is *beyond relational*, but to many people this is a confusing term. In one way of thinking, these are perfectly scalar types, but in yet another they really aren't.

The non-scalar types include the following:

- `hierarchyId`: Used to help build and manage a tree structure. It is very close to being a scalar type with several methods that can be applied to traverse and work with a hierarchy.
- Spatial types: geometry for dealing with planar/Euclidean (flat-Earth) data; geography for ellipsoidal (round-Earth) data, such as GPS longitude and latitude data. The spatial types technically hold arrays of values that represent sets on their own (and as you will see, you can join two shapes to see whether they overlap).
- XML: Used to store and manipulate XML values. A single XML column can more or less implement a database almost on its own.

One thing cannot be argued. Each of these types has some value to someone and fills a void that cannot be straightforwardly represented with the relational model, at least not as easily. I am a prude in many ways when it comes to normalization, but not every situation calls for strict adherence to the first normal form. What is, however, important is to know what you are doing and how you are violating the normal forms when you do and when it is appropriate.

So, are these bad things? My feeling on the matter is that it really depends upon usage. Since these types each have a lot more about them than simply what they can store, I have chosen to split them out into their own appendix. These types are mentioned here for completeness but will be covered in Appendix C (which will be available as a bonus download.)

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