

COL759 Assignment 2

Ashwini Kumar

2018MT60778

In this assignment we have implemented RSA Cryptosystem.

INPUT: Plaintext, a prime number 'r' and a number 'g' belonging to Z_p^* where $p=2*r+1$

OUTPUT:

User1(Server): Output for user 1 is the cipher text generated from the plaintext taken as input.

User2(Client): Output is the decrypted text received from user 1.

We find common session key(e) as following:

- (1) User 1 chooses a random key k_1 in Z_p^* and computes $x_1 = (g^{k_1}) \pmod p$
- (2) User 2 chooses a random key k_2 in Z_p^* and computes $x_2 = (g^{k_2}) \pmod p$
- (3) User 1 and User 2 exchange x_1 and x_2 .
- (4) User 1 computes $x_{1,2} = (x_2^{k_1}) \pmod p$
- (5) User 2 computes $x_{2,1} = (x_1^{k_2}) \pmod p$
- (6) $e=x_{1,2}$ if $x_{1,2}$ is odd and $e=x_{1,2}-1$ if $x_{1,2}$ is even

We find $e=x_{1,2}$ such that it has a modular inverse w.r.t mod $(p-1)$.

Then we calculate the private key as $d= (e^{-1}) \pmod{(p-1)}$.

User 1(Server):

Optimized Block size(b) is found by finding a value of b such that $(29^b) < p < (29^{b+1})$. This block size is used to break the plaintext into blocks each of size b. Then we calculate the value of M as:

Let $b=2$, first block is IN then $M= 13+ 8*29$

Then to find cipher text we calculate $C=(M^e) \pmod p = 26$ and then we write

$C=26+0*29$ and since $C < 29^b$, this results in $C=.AA$

User 2(Client):

User 2 receives the cipher text from User 1 and calculates value of C as done previously for calculating M. Then we calculate $M= (C^d) \pmod p$. Then we find plaintext from M as we did previously for C.

This way plaintext is divided into blocks by user 1 and then it sends the cipher text to user 2. Further user 2 receives the cipher text and then decrypts it to find the plaintext

HOW TO RUN:

Type

`python assignment2.py`

It will ask for taking r, g and plaintext as input so input

r

g

Plaintext

Output:

Two files User1.txt and User2.txt depicting some information on the values calculated