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Algorithm M
Input: A graph G = (V, E), IN, BN, LN, FL \subseteq V
Reduce G according to the reduction rules.
if there is some unreachable v \in \text{Free} \cup \text{FL} then return 0
if V = IN \cup LN then return |LN|
Choose a vertex v \in BN of maximum degree.
if d(v) > 3 or (d(v) = 2 and N_{EI}(v) \neq \emptyset) then
        \langle v \rightarrow \mathsf{LN} \mid \mid v \rightarrow \mathsf{IN} \rangle
                                                                                                                                                                            (B1)
else if d(v) = 2 then
       Let \{x_1, x_2\} = N_{\text{Free}}(v) such that d(x_1) < d(x_2).
       if d(x_1) = 2 then
              Let \{z\} = N(x_1) \setminus \{v\}
              if z \in Free then
                    \langle v \rightarrow \mathsf{LN} \mid \mid v \rightarrow \mathsf{IN}, x_1 \rightarrow \mathsf{IN} \mid \mid v \rightarrow \mathsf{IN}, x_1 \rightarrow \mathsf{LN} \rangle
                                                                                                                                                                            (B2)
              else if z \in FL then \langle v \to IN \rangle
       else if (N(x_1) \cap N(x_2)) \setminus FL = \{v\} and \forall z \in (N_{FL}(x_1) \cap N_{FL}(x_2)),
              d(z) > 3 then
                                                                                                                                                                            (B3)
              \langle v \rightarrow LN \mid \mid v \rightarrow IN, x_1 \rightarrow IN \mid \mid v \rightarrow IN, x_1 \rightarrow LN, x_2 \rightarrow IN \mid \mid
               v \to IN, x_1 \to LN, x_2 \to LN, N_{\text{Free}}(\{x_1, x_2\}) \to FL, N_{\text{RN}}(\{x_1, x_2\}) \setminus \{v\} \to LN
       else \langle v \to LN \mid v \to IN, x_1 \to IN \mid v \to IN, x_1 \to LN, x_2 \to IN \rangle
                                                                                                                                                                            (B4)
else if d(v) = 1 then
       Let P = (v = v_0, v_1, \dots, v_k) be a maximum path such that
              d(v_i) = 2, 1 \le i \le k, v_1, \dots, v_k \in Free.
       Let z \in N(v_{\nu}) \setminus V(P).
       if z \in FL and d(z) = 1 then \langle v_0, \dots, v_k \to IN, z \to LN \rangle
       else if z \in FL and d(z) > 1 then \langle v_0, \ldots, v_{k-1} \to IN, v_k \to LN \rangle
       else if z \in BN then \langle v \to LN \rangle
       else if z \in \text{Free then } \langle v_0, \dots, v_k \to \text{IN}, z \to \text{IN} \mid\mid v \to \text{LN} \rangle
                                                                                                                                                                            (B5)
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