Chapter 3

Processes

A system call is an interface between your program and the kernel. The linux kernel's job is to provide a variety of services to application programs and this is done using the provision of system calls.

3.1 Understanding Processes

The fundamental building block of each program is the process. A process is the name given to a program when it is loaded "for" the purpose of execution on the operating system. For a full description of what is comprised "inside" a process, please refer to your class-notes or slides.

To view a list of current processes in the system for a user, you may use the ps command.

ps

To view the complete list, you can adapt it to:

ps au

You will see plenty of columns as output. The columns you should be familiar with at this moment are underlined below:

- User: The owner of that process
- **PID**: The integer identifier
- CPU: Percent utilization of CPU
- <u>MEM</u>: Percent utilization of Memory space
- VSZ: Virtual Memory Size
- RSS: Non-swapped physical memory size

- TTY: Controlling Terminal
- <u>STAT</u>: The process states (D) Uninterruptible sleep (R) Running or ready (S) Interruptible sleep (T) Stopped (Z) Zombie (<) High Priority (N) Low Priority (s) Session leader, i.e., has child processes (l) multi-threaded (+) foreground
- START: When process has started
- TIME: Time running since
- <u>COMMAND</u>: The program/command used in this process.

The basic attribute of a process is it's ID (PID), and it's Parent ID (PPID). The system calls that can find the process ID, and the Parent Process ID are the getpid() and getppid() calls respectively. To use both of them, you will be required to include the sys/types.h and unistd.h C libraries.

The mere presence of the PPID means that there is a hierarchy of processes inside our operating system. Each process has a track of who are it's children, and each child process has a record of who is it's parent. The original process that is the grand-father of all processes is the *Init* Process. You can view this hierarchy using the pstree command.

pstree

This command will show a list of all processes currently in the system in the form of a tree.

3.1.1 Exercise

With respect to the myfirst.c file that you created in last labs, write a code that is able to find out the following:

- The PID value for myfirst.c
- The PPID value for myfirst.c
- The process name from the PPID value. (Note: You can use the system() call for this purpose. To find a process name from PPID, you would normally enter the following command on the shell, where 12345 is the ID of any process)

```
pstree -p | grep 12345
```

3.2 Process Lifecycle

Each process has to go through the following states during it's existence:

- 1. Creation
- 2. Running
- 3. Non-Running
 - (a) Ready
 - (b) Waiting
- 4. Termination

3.2.1 Creation States

About process creation concepts, understand and run the following code:

Q1 How many processes are created?

- Q2 Increase the value in for loop from i<1 to i<2 (i.e., 2 iterations in the loop). Compile and run your program. How many processes does it show this time? Draw a tree hierarchy of processes that you just created as given in Figure-3.1.
- Q3 Increase the value again to i<3 (i.e., 3 iterations). Compile and run your program. How many processes does it show? Draw a tree again. Why is it that we have called fork() 3 times in our code, yet we are seeing $2^n 1$ processes listed on screen?
- Q4 For fun, increase the value yet again to 100. Compile and run. What is going to happen? Does your OS Crash? Does your Program Crash? Can you modify your code to count the total number of fork()s made?

Try and run the following code:

```
01  #include <stdio.h>
02  #include <sys/types.h>
03  #include <unistd.h>
04  int main(void)
05  {
06  fork(); printf("He\n");
07  fork(); printf("Ha\n");
08  fork(); printf("Ho\n");
09  }
```

Note that we are calling each He, Ha, and Ho only once. Yet that does not appear when the program is run.

Q5 Can a Ho be output before a He? Why?

Explanation

The Fork() system call simply creates a new process which is exact replica of parent process. It requires the *unistd.h* C library. It's usage is as such:

```
// Fork code 1
#include <unistd.h>
#include <????> // See Question 2
int main()
{
   int p;
   p = fork();
   printf("Job Done\n");
}
```

To understand this code, try to answer the following questions:

Question1 We have used p = fork(). Why not simply fork()? Check **man fork** for answer.

Question 2 Check man page for printf. What library is used for this call?

Question3 Run your program. Why is it that printf() is used only once, yet we see the output "Job Done" displaying twice on our screen.

Question 4 Add the following statement to the end of your code and run it again. What output would you see?

```
printf("Value of P is %d\n", p);
```

So, the fork() system call will always return different kinds of values (As you should know from Question4). One of the values returned would be "0", the other will be a positive non-zero value.

The process that receives the "0" value is the child process.

The process that receives the positive non-zero value is the parent process.

Based on these values, we can then program our code in such a manner that both parent and child can do their own jobs and not appear to be doing the same things.

Why 0? Why Non-Zero?

The reason for this is that the child can always find out who it's parent is via the getp-pid() function call. However, it is difficult for the parent to know who it's children are. This can only be done if the value of the child's PID is returned to it at the time of the fork().

Look at the following code structure and implement it:

What would happen if we don't use the If/Else conditions and immediately write the two printf() statements?

Make sure you understand the structure of the Code, as well as what are the ways of knowing the following:

- 1. The PID
- 2. The Parent PID
- 3. The Child PID

3.2.1.1 Exercise 1

Modify the fork() call code 2 above and add a system call for sleep() after both the printf() statements. Give a time of 120 seconds to the sleep call. While both the parent and child are sleeping, open a new terminal and check out the outcome of **pstree** command as well as the **ps** command. Study the output for your parent and child process in both commands, especially noting the STAT column.

3.2.1.2 Exercise 2

Model a fork() call in C/C++ so that your program can create a total of EXACTLY 6 processes (including the parent). (Note: You may check the number of processes created using the method from Exercise 1 (Section-3.2.1.1 by using a sleep of 60 seconds or more and entering eithe ps, or pstree in another terminal). Your process hierarchy should be as follows:

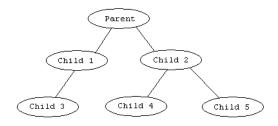


Figure 3.1: Exercise 2 Model

3.2.2 Running States

The job of the exec() call is to <u>replace</u> the current process with a new process/program. Once the exec() call is made, the current process is "gone" and a new process starts. The Exec() call is actually a family of 6 system calls. Difference between all 6 can be seen from man 3 exec

- int execl(const char *path, const char *arg, ...);
- int execlp(const char *file, const char *arg, ...);
- int execle(const char *path, const char *arg, char *const envp[]);
- int execv(const char *path, char *const argv[]);
- int execve(const char *path, const char *argv[], char *const envp[]);
- int execvp(const char *file, char *const argv[]);

Following is usage of one flavor of Exec, the execv() system call:

```
#include <unistd.h>
#include <stdio.h>
int main()
{
   int p;
   char *arg[] = {"/usr/bin/ls", 0};
   p = fork();
   if (p == 0)
   {
      printf("Child Process\n");
      execv(arg[0], arg);
```

```
printf("Child Process\n");
}
else
{
   printf("Parent Process\n");
}
```

Again look at the code. We have referred to Exec() system call but in code we see Execv(). Read man pages for this and find out. To help you understand, try finding answers to the following:

Question 1 What is the 1st argument to the execv() call? What is it's contents?

Question 2 What is the 2nd argument to it? What is it's contents?

Question 3 What is arg?

Question4 Look at the code of the child process (p==0). How many times does the statement "Child Process" appear? Why?

Diagrammatically, the functioning of the exec() system calls would be represented in Figure 3.2, where the dotted line mark the execution and transfer of control whereas the straight lines mark the waiting time.

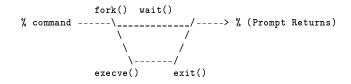


Figure 3.2: Fork() Exec() representation

3.2.3 Waiting States

3.2.3.1 Sleep()

The sleep() call can be used to cause delay in execution of a program. The delay can be provided as an integer number (representing number of seconds). Usage is simply sleep(int), which will delay the process execution for int seconds. To use sleep(), you will require the unistd.h C library.

Exercise

Write a C program that can display a count from 10 to 0 (reverse order) using a for or a while loop. Each number should be displayed after delay of 1 second.

Termination States 3.2.4

A process can terminate in one of the following ways:

- 1. Main function calls return
- 2. Exit function called
- 3. Aborted by higher priority process (e.g., parent or kernel)
- 4. Terminated by a signal

Regardless of any of these reasons, the same kernel code is invoked which performs the following:

- 1. Close open files
- 2. Notifies Parent and Children
- 3. Release memory resources

3.2.4.1Exit() Call

void anotherExit()

The exit() system call causes a normal program to terminate and return status to the parent process. Study the behaviour of the following program. You will see that there may be a number of exit points from a program. Can you identify which exit() call is being used each time a program exits??

```
#include <stdio.h>
int main()
  int num;
  void anotherExit(); // Function Prototypeanswer the questions in the end.
 printf("Enter a Number: ");
  scanf("%i", num);
  if (num>25)
   printf("exit 1\n");
    exit(1);
  }
  else
    anotherExit();
```

```
printf("Exit 2\n");
exit(1);
```

3.2.4.2Atexit() Call

The code below provides two functions; atexit() and exit(). Note the structure of code and the behaviour of execution and then attempt the questions in the end.

```
01 #include<stdio.h>
02 #include<stdlib.h>
   int main(void) {
      void f1(void), f2(void), f3(void);
04
05
      atexit(f1);
06
      atexit(f2);
      atexit(f3):
07
08
     printf("Getting ready to exit\n");
09
      exit(0):
10 }
11
   void f1(void) {
     printf("In f1\n");
12
13 }
14
   void f2(void) {
15
     printf("In f2\n");
16
    void f3(void) {
17
18
     printf("In f3\n");
19
```

- Q1 What is the difference between exit() and atexit()? What do they do? (Check man atexit and man 3 exit).
- Q2 What does the 0 provided in the exit() call mean? What will happen if we change it to 1? (Check manual page for exit)
- Q3 If we add an exit call to function f1, f2, or f3. What will happen to execution of our program?
- Q4 Why do you think we are getting reverse order of execution of atexit calls?

3.2.4.3Abort Call

Type, run and execute the code below. Then

```
01 #include <stdio.h>
   #include <stdlib.h>
03 int main() {
04
     abort():
05
      exit(0);
```

Q1 Check the man pages for abort. does the abort call terminate the program? What is the name of the particular signal?

- **Q2** Execute your program. What is the output of our program?
- Q3 Include the abort call in function f3 in our code provided for Atexit() call. How does our program terminate using this?

3.2.4.4 Kill Call

Full description of this call will appear in Section-5.1.2. But as a demonstration, you may run the following code to see how the kill() call works.

```
#include <stdio.h>
#include <sys/types.h>
#include <signal.h>
int main()
{
  printf("Hello");
  kill(getpid(), 9);
  printf("Goodbye");
}
```

You are already familiar with getpid() system call. To find out what 9 is, first look at the output of the command:

```
kill -1
```

Find out the word mentioned next to 9. Search on the internet for this.

3.3 Death of Parent or Child

We covered the process termination using the exit() system call and that a program may have more than one exit points. In a relationwhip where a process has spawned children, what would be the effect if either the parent dies or the child dies on the other processes linked to it? This section will look and study such a development.

3.3.1 Parent Dies Before Child

If a parent process dies before it's children, the children will be orphaned. They will be assigned to another process in the system and as such the children should be informed about who their new parent is going to be. This new parent process is the parent of all processes in the system, i.e., the "init" process.

```
#include <unistd.h>
#include <stdio.h>
int main()
  int i, pid;
  pid = fork();
  if (pid > 0)
                        // Parent
    sleep(2);
    exit(0);
  }
  else if (pid == 0)
                        // Child
    for (i=0; i < 5; i++)
      printf("My parent is %d\n", getppid());
      sleep(1);
 }
}
```

Run the code.

- 1. What are the PPID values you are receiving from the for loop?
- 2. What has happened when the numbers of the PPID change?
- 3. What is now PID of the init process?

3.3.2 Child Dies Before Parent

Following code is complete opposite of what we **Or** have seen so far.

Run the code. In another window (terminal), check the status of your parent and child processes using the "ps" command. You should be able to see a Z, or <defunct> next to the child process which has been created. Such a process is usually termed a "Zombie" process.

Wait for 60 seconds, then check the status of your process again. You will note that both the parent, as well as the Zombie process are now gone.