

# Chapter 6

## IPC: Pipes

Pipes are another IPC technique for processes to communicate with one another. It can also be used by two threads within the same process to communicate.

In the description of file descriptors from Chapter-4, you will recall that there are three files that are open all the time for input and output purposes. These are:

1. The standard input
2. The standard output
3. The standard error

And that they have a file-descriptor of 0, 1, and 2 respectively. Whenever a program needs to display some output (via `cout` or `printf`), it will write that output to the standard output file descriptor. This will in return be displayed on the monitor. Whenever a program needs to take some input from the keyboard, it will take it's input from the standard input file descriptor. Similarly, whenever an error needs to be displayed, that error will be sent to the standard error file descriptor. These files are linked-up internally to peripheral devices such as keyboard, monitor, etc. However, these linkages can be changed to point to something else. Pipes work by doing exactly that! So a pipe will::

- Redirect the standard output of one process to become the standard input of another

The rest of communication is done using the following rules:

- The pipe will be a buffer region in main memory which will be accessible by only two processes.

- One process will read from the buffer while the other will write to it.
- One process cannot read from the buffer unless and until the other has written to it.

### 6.1 Pipe On the Shell

Run the following command:

```
pstree
```

As you will notice, the output is too long to fit in the screen. Now run the command with:

```
pstree | less
```

Using the up and down arrows you will notice that you can browse through the output which was otherwise not visible in the first command which we gave. To exit, press **q**.

The `|` is the symbol for pipe and as you would have guessed, `pstree` and `less` are two processes. In this usage, the *standard output of pstree* has become the *standard input of the less* program.

Try the following command:

```
pstree | grep bash
```

Again, you will see that the output of above command is much different from just `pstree` command. What the above command should print is only those lines of text from the `pstree` output in which the keyword `bash` appears. Hence, `pstree` and `grep` are two separate processes. But the standard output of the `pstree` command has become the standard input of the `grep` program. The `pstree` command writes out its output to the `grep` process. The `grep` process receives it, searches for keyword, formats output, and then displays the result.

## 6.2 Pipe System Call

Type, compile, and run the following code:

```
#include <unistd.h>
int main()
{
    int pfd[2];
    pipe(pfd);
}
```

The above code creates a pipe using the `pipe()` system call. We have passed it the name of the integer array `pfd` that we have declared earlier on.

**Task:** Rewrite this code so that you can view the contents of the array using `printf` arguments. You should see two numbers. What are these numbers?

Let us extend our code. Type, compile, and run the following:

```
01 #include <unistd.h>
02 int main()
03 {
04     int pid;           // for storing fork() return
05     int pfd[2];        // for pipe file descriptors
06     char aString[20]; // Temporary storage
07     pipe(pfd);         // create our pipe
```

When line number 07 completes, we will have the following in our process:

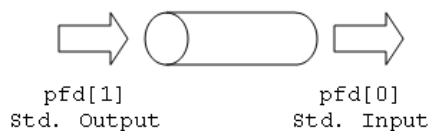


Figure 6.1: Before `Fork()`

```
08     pid = fork();      // create child process
```

When line number 08 completes, we will have the following in our two processes:

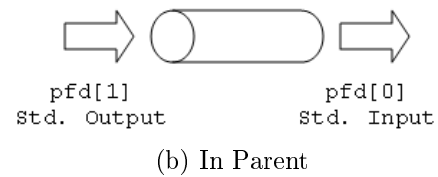
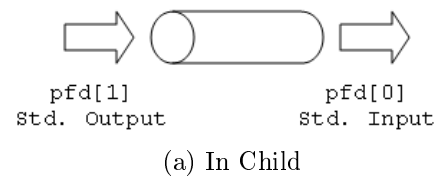


Figure 6.2: After `Fork()` is made

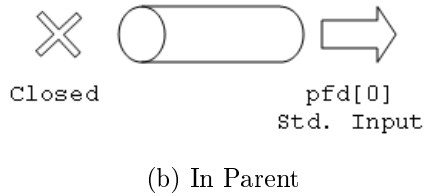
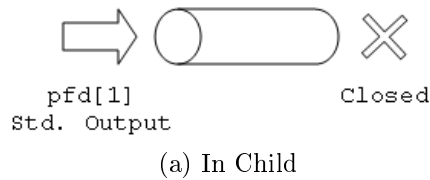
Continuing with rest of code:

```
09     if (pid == 0)      // For child
10     {
11         write(pfd[1], "Hello", 5); // Write onto pipe
12     }
13     else                // For parent
14     {
15         read(pfd[0], aString, 5);  // Read from pipe
16     }
17 }
```

Just like we open a file, we read from a file, we write to a file, and we close a file, we will perform the same operations of `open()`, `close()`, `read()` and `write()` on the pipe.

**Task:** Rewrite this code so that you can see the contents of `aString` in the parent before and after the `read()` call. What are the contents? You will notice that “Hello” has been mentioned in the child process. Then how is it possible that we are able to see the term “Hello” in the parent process? The answer is through the pipe mechanism which we just used.

In the code we just saw, since the child is only going to write to a pipe and the parent will only read from the pipe, it makes sense to close the read capabilities for the child and write capabilities for the parent for that particular pipe. Diagrammatically, we want to achieve something like the following:



```

        close(pfd[1]);
        execlp("ls", "ls", (char *) 0);
    }
}

```

Figure 6.3: Closing un-necessary ends of the Pipe

For that, we have to add the following before line 11:

```
close(pfd[0]);
```

and the following before line 15:

```
close(pfd[1]);
```

### 6.2.1 Example

Type, run and execute the code below. It should give output which is equivalent to the command *ls / wc* (Wc is used for printing three numbers; the number of newlines, the number of words, and the byte count for a file). Note the usage of pipes.

```

#include <unistd.h>
#include <string.h>
#include <stdio.h>
int main()
{
    int pfd[2];
    pipe(pfd);
    if (fork() == 0)
    {
        close(pfd[1]);
        dup2(pfd[0], 0);
        close(pfd[0]);
        execlp("wc", "wc", (char *) 0);
    }
    else
    {
        close(pfd[0]);
        dup2(pfd[1], 1);
    }
}

```