7 Data Structure

7.1 Link-Cut Tree . . . . . . . . . . . . . . . . .

#### Contents 8 Others 25 8.1 Find max tangent(x,y is increasing) . . . . . . 25 25 1 Basic 1.1 .vimrc 8.4 Hilbert Curve . . . . 1.2 Custom Hash . . . . . . . . . . . . . . . . . 1.3 python-related . . . . . . . . . 1 Basic 2 flow 1.1 .vimrc 2.1 ISAP . syn on se ai nu rnu ru cul mouse=a se cin et ts=2 sw=2 sts=2 so \$VIMRUNTIME/mswin.vim colo desert filet plugin indent on 2.9 Max flow with lower/upper bound . . . . . . . . . . . no <F5> :!./a.out<CR> no <F9> :!g++ -02 -std=gnu++14 -lm % -g -fsanitize= 2.11Flow Method . . . . . . . . . . . . . . . undefined -Wall -Wextra -Wshadow -Wno-unused-result <CR> 3 Math 3.1 FFT . . . . . . . . . . . . . . . . . 1.2 Custom Hash 3.3 Fast Walsh Transform . . . . 3.4 Poly operator . . . . . . struct custom\_hash { static uint64\_t splitmix64(uint64\_t x) { 3.6 BerlekampMassey . . . . . . . . . . . . . . . . x += 0x9e3779b97f4a7c15;3.7 Miller Rabin . . . . . . . . . . . . . . . . . . $x = (x \land (x >> 30)) * 0xbf58476d1ce4e5b9;$ $x = (x \land (x >> 27)) * 0x94d049bb133111eb;$ 3.8 Simplex . . . . . . . . . . . . . . . . . 3.10Chinese Remainder . . . . . . . . . . . . . . . . return $x \wedge (x \gg 31)$ ; size\_t operator()(uint64\_t x) const { 3.13Discrete sqrt . . . . . . . . . . . . . . . . . static const uint64\_t FIXED\_RANDOM = chrono:: steady\_clock::now().time\_since\_epoch().count(); return splitmix64(x + FIXED\_RANDOM); }; 1.3 python-related 4.1 Intersection of 2 lines . . . . . . . . . . . . . . . . . . 10 from fractions import Fraction 4.3 Intersection of 2 segments . . . . . . . . . . . . . . . . 10 from decimal import Decimal, getcontext 4.4 Banana . . . 10 4.5 Intersection of circle and line $\dots \dots$ getcontext().prec = 250 # set precision 4.6 Intersection of polygon and circle . . . . . . . . 10 4.7 Intersection of 2 circles . . . . . . . . . . . . 10 itwo = Decimal(0.5)10 two = Decimal(2)4.9 Li Chao Segment Tree ........ 11 11 N = 2004.11 Tangent line of two circles . . . . . . . . . . . . . def angle(cosT): 4.12 Tangent line of point and circle . . . . . . . . . . . 12 """given cos(theta) in decimal return theta""" 4.13Min distance of two convex . . . . . . . . . . . . 12 12 for i in range(N): cosT = ((cosT + 1) / two) \*\* itwo13 sinT = (1 - cosT \* cosT) \*\* itwo return sinT \* (2 \*\* N) 13 4.18Min Enclosing Ball . . . . . . . . . . . . . . . . . 14 pi = angle(Decimal(-1))15 4.20Min dist on Cuboid $\dots$ ...... 15 2 flow 5 Graph 2.1 ISAP #define SZ(c) ((int)(c).size()) struct Maxflow { 5.4 Strongly Connected Component . . . . . . . . . . . . . . . . . . 16 static const int MAXV = 20010; 5.6 Maximum General graph Matching . . . . . . . . . . . . . . . . static const int INF = 1000000; 17 5.7 Minimum General Weighted Matching . . . . . . . . . . 17 struct Edge { 5.8 Maximum General Weighted Matching . . . . . . . . 18 int v, c, r; 19 Edge(int \_v, int \_c, int \_r): v(\_v), c(\_c), r(\_r) {} 5.10BCC based on vertex . . . . . . . . . . . 19 5.12Directed Graph Min Cost Cycle . . . . . 20 int s, t; 5.13K-th Shortest Path . . . . . . . . . . . . 21 vector<Edge> G[MAXV]; 5.14Chordal Graph . . . . . . . . . . . . . 22 int iter[MAXV], d[MAXV], gap[MAXV], tot; void init(int x) { 5.15Graph Method . . . . . . . . 6 String tot = x+2; 22 s = x+1, t = x+2; for(int i = 0; i <= tot; i++) { G[i].clear(); 23 iter[i] = d[i] = gap[i] = 0;void addEdge(int u, int v, int c) { 6.9 Cyclic LCS . . . . . . . . . . . . . . . G[u].push\_back(Edge(v, ć, SZ(G[v]) )); G[v].push\_back(Edge(u, 0, SZ(G[u]) - 1));

```
National Taiwan University CRyptoGRapheR
  int dfs(int p, int flow) {
    if(p == t) return flow;
    for(int &i = iter[p]; i < SZ(G[p]); i++) {</pre>
      Edge \&e = G[p][i]
      if(e.c > 0 \& d[p] == d[e.v]+1)
        int f = dfs(e.v, min(flow, e.c));
        if(f) {
          G[e.v][e.r].c += f;
          return f;
      }
    if( (--gap[d[p]]) == 0) d[s] = tot;
    else {
      d[p]++;
      iter[p] = 0;
      ++gap[d[p]];
    return 0;
  int solve() {
    int_res = 0;
    gap[0] = tot;
    for(res = 0; d[s] < tot; res += dfs(s, INF));
    return res;
} flow;
2.2 MinCostFlow
struct zkwflow{
  struct Edge {
    int to, rev, cap; ll cost;
  vector<Edge> g[N];
```

```
int nv, sv, tv, ptr[N];
bool vis[N]; ll dist[N];
void init(int n,int s,int t){
  nv=n+1; sv=s; tv=t;
for(int i=0;i<nv;i++) g[i].clear();</pre>
void add_edge(int a, int b, int c, ll w) {
  g[a].push_back(Edge{b,int(g[b].size()),c,w});
  g[b].push_back(Edge\{a,int(g[a].size())-1,0,-w\});
bool augment() { // SPFA
  for (int i = 0; i < nv; i++) {
     dist[i] = LLÍNF; vis[i] = false;
  dist[sv] = 0;
  vector<int> que = { sv };
for (int i = 0; i < int(que.size()); i++) {</pre>
     int_v_= que[i];
     vis[v] = false;
     for (auto& e : g[v]) {
  if (e.cap == 0 || dist[e.to] <= dist[v] + e.</pre>
            cost)
          continue;
       dist[e.to] = dist[v] + e.cost;
       if (!vis[e.to]) {
          vis[e.to] = true
          que.push_back(e.to);
    }
  }
  return dist[tv] != LLINF;
int dfs(int v, int r) {
  if (v == tv) return r;
  vis[v] = true;
  int res = 0;
  for (int& i = ptr[v]; i < int(g[v].size()); i++) {</pre>
     Edge& e = g[v][i];
if (e.cap == 0 || dist[e.to] != dist[v] + e.cost
          | vis[e.to])
       continue
     int d = dfs(e.to, min(r - res, e.cap));
     res += d; e.cap -= d;
     g[e.to][e.rev].cap += d;
     if (res == r) {
       vis[v] = false;
```

```
break;
    }
}
return res;
}
pair<int, ll> solve() {
    int flow = 0; ll cost = 0;
    while (augment()) {
        fill_n(ptr, nv, 0);
        int d = dfs(sv, INF);
        flow += d; cost += d * dist[tv];
    }
    return { flow, cost };
}
}flow;
```

#### 2.3 Dinic

```
struct Dinic{
  static const int MXN = 10000;
struct Edge{ int v,f,re; };
   int n,s,t,level[MXN];
   vector<Edge> E[MXN];
  void init(int _n, int _s, int _t){
  n = _n;  s = _s;  t = _t;
  for (int i=0; i<n; i++) E[i].clear();</pre>
   void add_edge(int u, int v, int f){
     E[u].PB({v,f,(int)E[v].size()})
     E[v].PB({u,0,(int)E[u].size()-1});
   bool BFS(){
     for (int i=0; i<n; i++) level[i] = -1;</pre>
     queue<int> que;
     que.push(s):
     level[s] = 0;
     while (!que.empty()){
        int u = que.front(); que.pop();
        for (auto &it : E[u]){
            f (it.f > 0 && level[it.v] == -1){
level[it.v] = level[u]+1;
            que.push(it.v);
       }
     return level[t] != -1;
   int DFS(int u, int nf){
     if (u == t) return nf;
     int res = 0;
     for (auto &it : E[u]){
       if (it.f > 0 && level[it.v] == level[u]+1){}
          int tf = DFS(it.v, min(nf,it.f));
          res += tf; nf -= tf; it.f -= tf;
          E[it.v][it.re].f += tf;
          if (nf == 0) return res;
       }
     if (!res) level[u] = -1;
     return res;
   int flow(int res=0){
     while ( BFS() )
       res += DFS(s,2147483647);
     return res;
}flow;
```

# 2.4 Kuhn Munkres

```
struct KM{ // max weight, for min negate the weights
   static const int MXN = 2001; // 1-based
   static const ll INF = 0x3f3f3f3f;
   int n, mx[MXN], my[MXN], pa[MXN];
   ll g[MXN][MXN], lx[MXN], ly[MXN], sy[MXN];
   bool vx[MXN], vy[MXN];
   void init(int _n) {
        n = _n;
        for(int i=1; i<=n; i++) fill(g[i], g[i]+n+1, 0);
   }
   void addEdge(int x, int y, ll w) {g[x][y] = w;}
   void augment(int y) {</pre>
```

```
for(int x, z; y; y = z)
    x=pa[y], z=mx[x], my[y]=x, mx[x]=y;
   void bfs(int st) {
     for(int i=1; i<=n; ++i) sy[i]=INF, vx[i]=vy[i]=0;</pre>
     queue<int> q; q.push(st);
     for(;;) {
        while(q.size()) {
          int x=q.front(); q.pop(); vx[x]=1;
for(int y=1; y<=n; ++y) if(!vy[y]){
    ll t = lx[x]+ly[y]-g[x][y];</pre>
             if(t==0){
                pa[y]=x
                if(!my[y]){augment(y);return;}
                vy[y]=1, q.push(my[y]);
             }else if(sy[y]>t) pa[y]=x,sy[y]=t;
        11 cut = INF;
        for(int y=1; y<=n; ++y)</pre>
          if(!vy[y]&&cut>sy[y]) cut=sy[y];
        for(int j=1; j<=n; ++j){
  if(vx[j]) lx[j] -= cut;
  if(vy[j]) ly[j] += cut;</pre>
          else sy[j] -= cut;
        for(int y=1; y<=n; ++y) if(!vy[y]&&sy[y]==0){
  if(!my[y]){augment(y);return;}</pre>
          vy[y]=1, q.push(my[y]);
     }
  11 solve(){
     fill(mx, mx+n+1, 0); fill(my, my+n+1, 0);
     fill(ly, ly+n+1, 0); fill(lx, lx+n+1, -INF);
     for(int x=1; x<=n; ++x) for(int y=1; y<=n; ++y) lx[x] = max(lx[x], g[x][y]);
     for(int x=1; x<=n; ++x) bfs(x);</pre>
     11 \text{ ans} = 0;
     for(int y=1; y<=n; ++y) ans += g[my[y]][y];
     return ans;
}graph;
```

## 2.5 DMST

```
* Edmond's algoirthm for Directed MST
   runs in O(VE)
const int MAXV = 10010;
const int MAXE = 10010
const int INF = 2147483647;
struct Edge{
  int u, v, c;
  Edge(int x=0, int y=0, int z=0) : u(x), v(y), c(z){}
int V, E, root
Edge edges[MAXE];
inline int newV(){ return ++ V; }
inline void addEdge(int u, int v, int c)
{ edges[++E] = Edge(u, v, c); }
bool con[MAXV];
int mnInW[MAXV], prv[MAXV], cyc[MAXV], vis[MAXV];
inline int DMST(){
  fill(con, con+V+1, 0);
  int r1 = 0, r2 = 0;
  while(1){
    fill(mnInW, mnInW+V+1, INF);
    fill(prv, prv+V+1, -1);
REP(i, 1, E){___
       int u=edges[i].u, v=edges[i].v, c=edges[i].c;
      if(u != v && v != root && c < mnInW[v])
        mnInW[v] = c, prv[v] = u;
    fill(vis, vis+V+1, -1);
    fill(cyc, cyc+V+1, -1);
    r1 = 0;
bool jf = 0;
    REP(i, 1, V){
      if(con[i]) continue ;
if(prv[i] == -1 && i != root) return -1;
```

```
if(prv[i] > 0) r1 += mnInW[i];
      for(s = i; s != -1 && vis[s] == -1; s = prv[s])
        vis[s] = i;
      if(s > 0 \& vis[s] == i){
         // get a cycle
        jf = 1; int v = s;
        do{
          cyc[v] = s, con[v] = 1;
          r2 += mnInW[v]; v = prv[v];
        }while(v != s);
        con[s] = 0;
      }
    if(!jf) break ;
    REP(i, 1, E){
      int &u = edges[i].u;
int &v = edges[i].v;
      if(cyc[v] > 0) edges[i].c -= mnInW[edges[i].v];
      if(cyc[u] > 0) edges[i].u = cyc[edges[i].u];
      if(cyc[v] > 0) edges[i].v = cyc[edges[i].v];
      if(u == v) edges[i--] = edges[E--];
  return r1+r2;
}
```

#### SW min-cut 2.6

```
const int INF=0x3f3f3f3f3f;
template<typename T>
struct stoer_wagner{// 0-base
  static const int MAXN=501;
  T g[MAXN] [MAXN], dis[MAXN];
  int nd[MAXN],n,s,t;
  void init(int _n){
     for(int i=0;i<n;++i)</pre>
       for(int j=0;j<n;++j)g[i][j]=0;</pre>
  void add_edge(int u,int v,T w){
    g[u][v]=g[v][u]+=w;
  T min_cut(){
    T ans=INF;
     for(int i=0;i<n;++i)nd[i]=i;</pre>
     for(int ind,tn=n;tn>1;--tn){
       for(int i=1;i<tn;++i)dis[nd[i]]=0;</pre>
       for(int i=1;i<tn;++i){</pre>
         ind=i;
         for(int j=i;j<tn;++j){
  dis[nd[j]]+=g[nd[i-1]][nd[j]];</pre>
           if(dis[nd[ind]]<dis[nd[j]])ind=j;</pre>
         swap(nd[ind],nd[i]);
       if(ans>dis[nd[ind]])
         ans=dis[t=nd[ind]],s=nd[ind-1];
       for(int i=0;i<tn;++i)</pre>
         g[nd[ind-1]][nd[i]]=g[nd[i]][nd[ind-1]]
                                +=g[nd[i]][nd[ind]];
    return ans:
  }
};
```

#### Max Cost Circulation

```
struct MaxCostCirc {
     static const int MAXN = 33;
    int n , m;
struct Edge { int v , w , c , r; };
vector<Edge> g[ MAXN ];
int dis[ MAXN ] , prv[ MAXN ] , prve[ MAXN ];
bool vis[ MAXN ];
     int ans;
    void init( int _n , int _m ) : n(_n), m(_m) {}
void adde( int u , int v , int w , int c ) {
   g[ u ].push_back( { v , w , c , SZ( g[ v ] ) } );
   g[ v ].push_back( { u , -w , 0 , SZ( g[ u ] )-1 } )
```

```
bool poscyc() {
       fill( dis , dis+n+1 , 0 );
       fill( prv , prv+n+1 , 0 );
       fill( vis , vis+n+1 , 0 );
        int tmp = -1;
       Int tmp = -1;
FOR( t , n+1 ) {
    REP( i , 1 , n ) {
    FOR( j , SZ( g[ i ] ) ) {
        Edge& e = g[ i ][ j ];
        if( e.c && dis[ e.v ] < dis[ i ]+e.w ) {
            dis[ e.v ] = dis[ i ]+e.w;
            prv[ e.v ] = i;
            prv[ e.v ] = i;
            prv[ e.v ] = i;
            prv[ e.v ] = i;</pre>
                     prve[ e.v ] = j;
                     if(\bar{t} = n)
                        tmp = i;
                        break;
       int cur = tmp;
       while( !vis[ cur ] ) {
  vis[ cur ] = 1;
  cur = prv[ cur ];
        int now = cur, cost = 0, df = 100000;
       qo{
          Edge &e = g[ prv[ now ] ][ prve[ now ] ];
df = min( df , e.c );
           cost += e.w;
           now = prv[ now ];
       }while( now != cur );
       ans += df*cost; now = cur;
          Edge &e = g[ prv[ now ] ][ prve[ now ] ];
Edge &re = g[ now ][ e.r ];
           e.c -= df;
           re.c += df;
       now = prv[ now ];
}while( now != cur );
       return 1;
} circ;
```

#### 2.8 Gomory-Hu Tree

```
//n,Dinic::flow must be filled
//result:e[u][v]=u-v mincut;p[u]:u's parent on cut tree
int n,e[MXN][MXN],p[MXN];
void gomory_hu(){
  fill(p, p+n, 0);
fill(e[0], e[n], INF);
  for(int s = 1; s < n; s++){
    int t = p[s];
Dinic F; F.init(n,s,t);
    copy(flow.E,flow.E+MXN,F.E);
    if ( p[i] == t && F.level[i]!=-1 ) p[i] = s;
}
```

## 2.9 Max flow with lower/upper bound

```
// Max flow with lower/upper bound on edges
// use with ISAP
int in[ N ] , out[ N ];
int l[ M ] , r[ M ] , a[ M ] , b[ M ];
int solve(int n, int m, int s, int t){
     flow.init( n );
    frow.trt( n ),
for( int i = 0 ; i < m ; i ++ ){
  in[ r[ i ] ] += a[ i ];
  out[ l[ i ] ] += a[ i ];
  flow.addEdge( l[ i ] , r[ i ] , b[ i ] - a[ i ] );
  // flow from l[i] to r[i] must in [a[ i ], b[ i ]]</pre>
     int nd = 0;
    for( int i = 0 ; i <= n ; i ++ ){
  if( in[ i ] < out[ i ] ){
    flow.addEdge( i , flow.t , out[ i ] - in[ i ] );
    nd += out[ i ] - in[ i ];
}</pre>
```

```
if( out[ i ] < in[ i ] )</pre>
       flow.addEdge( flow.s , i , in[ i ] - out[ i ] );
  // original sink to source
  flow.addEdge( t , s , INF );
if( flow.solve() != nd )
    // no solution
    return -1;
  int ans = flow.G[ s ].back().c; // source to sink
  flow.G[s].back().c = flow.G[t].back().c = 0;
  // take out super source and super sink
  for( size_t i = 0 ; i < flow.G[ flow.s ].size() ; i</pre>
    flow.G[flow.s][i].c = 0;
    Maxflow::Edge &e = flow.G[ flow.s ][ i ];
    flow.G[ e.v ][ e.r ].c = 0;
  for( size_t i = 0 ; i < flow.G[ flow.t ].size() ; i</pre>
    flow.G[ flow.t ][ i ].c = 0;
    Maxflow::Edge &e = flow.G[ flow.t ][ i ];
    flow.G[ e.v ][ e.r ].c = 0;
  flow.addEdge( flow.s , s , INF );
flow.addEdge( t , flow.t , INF );
flow.reset(); // set iter,d,gap to 0
  return ans + flow.solve();
2.10 HLPPA
```

```
template <int MAXN, class T = int>
struct HLPP {
  const T INF = numeric_limits<T>::max();
  struct Edge {
    int to, rev; T f;
  int n, s, t;
  vector<Edge> adj[MAXN];
  deque<int> lst[MAXN];
vector<int> gap[MAXN];
  int ptr[MAXN];
  T ef[MAXN];
  int h[MAXN], cnt[MAXN], work, hst=0/*highest*/;
  void init(int _n, int _s, int _t) {
    n=_n+1;    s = _s;    t = _t;
    for(int i=0;i<n;i++) adj[i].clear();</pre>
  void addEdge(int u,int v,T f,bool isDir = true){
  adj[u].push_back({v,adj[v].size(),f});
    adj[v].push_back({u,adj[u].size()-1,isDir?0:f});
  void updHeight(int v, int nh) {
    work++;
    if(h[v] != n) cnt[h[v]]--;
    h[v] = nh;
    if(nh == n) return;
    cnt[nh]++, hst = nh; gap[nh].push_back(v);
    if(ef[v]>0) lst[nh].push_back(v), ptr[nh]++;
  void globalRelabel() {
    work = 0;
    fill(h, h+n, n);
     fill(cnt, cnt+n, 0);
     for(int i=0; i<=hst; i++)</pre>
    lst[i].clear(), gap[i].clear(), ptr[i] = 0;
queue<int> q({t}); h[t] = 0;
    while(!q.empty()) {
  int v = q.front(); q.pop();
       for(auto &e : adj[v])
         if(h[e.to] == n \& adj[e.to][e.rev].f > 0)
            q.push(e.to), updHeight(e.to, h[v] + 1);
       hst = h[v];
    }
  void push(int v, Edge &e) {
    if(ef[e.to] == 0)
       lst[h[e.to]].push_back(e.to), ptr[h[e.to]]++;
    T df = min(ef[v], e.f);
    e.f -= df, adj[e.to][e.rev].f += df;
    ef[v] -= df, ef[e.to] += df;
```

```
void discharge(int v) {
    int nh = n;
     for(auto &e : adj[v]) {
       if(e.f > 0) {
         if(h[v] == h[e.to] + 1) {
           push(v, e);
           if(ef[v] <= 0) return;</pre>
         else nh = min(nh, h[e.to] + 1);
       }
    if(cnt[h[v]] > 1) updHeight(v, nh);
       for(int i = h[v]; i < n; i++) {
         for(auto j : gap[i]) updHeight(j, n);
gap[i].clear(), ptr[i] = 0;
    }
  T solve() {
    fill(ef, ef+n, 0);
ef[s] = INF, ef[t] = -INF;
    globalRelabel();
     for(auto &e : adj[s]) push(s, e);
    for(; hst >= 0; hst--)
       while(!lst[hst].empty())
         int v=lst[hst].back(); lst[hst].pop_back();
         discharge(v);
if(work > 4 * n) globalRelabel();
    return ef[t] + INF;
};
```

#### 2.11 Flow Method

```
Maximize c^T x subject to Ax \le b, x \ge 0;
with the corresponding symmetric dual problem,
Minimize b^T y subject to A^T y \geq c, y \geq 0.
Maximize c^T x subject to Ax \le b;
with the corresponding asymmetric dual problem,
Minimize b^T y subject to A^T y = c, y \geq 0.
General Graph:
|Max Ind. Set| + |Min Vertex Cover| = |V|
|Max Ind. Edge Set| + |Min Edge Cover| = |V|
Bipartite Graph:
|Max Ind. Set| = |Min Edge Cover|
| Max Ind. Edge Set| = | Min Vertex Cover|
To reconstruct the minimum vertex cover, dfs from each
```

unmatched vertex on the left side and with unused edges only. Equivalently, dfs from source with unused edges only and without visiting sink. Then, a vertex is chosen iff. it is on the left side and without visited or on the right side and visited through dfs.

Minimum Weighted Bipartite Edge Cover:

Construct new bipartite graph with n+m vertices on each side:

for each vertex u, duplicate a vertex u' on the other side

for each edge (u,v,w), add edges (u,v,w) and (v',u',w)for each vertex u, add edge (u,u',2w) where w is min edge connects to u

then the answer is the minimum perfect matching of the new graph (KM)

Maximum density subgraph ( \sum{W\_e}+ \sum{W\_v} ) / |V| Binary search on answer:

For a fixed D, construct a Max flow model as follow: Let S be Sum of all weight( or inf)

1. from source to each node with cap = 9

- 2. For each (u,v,w) in E, (u->v,cap=w), (v->u,cap=w)
- 3. For each node v, from v to sink with cap = S + 2 \* D deg[v] 2 \* (W of v)

where  $deg[v] = \sum weight of edge associated with v$ If maxflow < S \* IVI, D is an answer.

```
Requiring subgraph: all vertex can be reached from
    source with
edge whose cap > 0.
```

Maximum closed subgraph

- connect source with positive weighted vertex( capacity=weight)
- connect sink with negitive weighted vertex(capacity =-weiaht)
- 3. make capacity of the original edges = inf
- 4. ans = sum(positive weighted vertex weight) (max flow)

Minimum Path Cover of DAG

- 1. For each vertex v, split it to v\_in and v\_out.
- 2. For each edge (u->v), add an edge between u\_out and v in
- 3. |Minimum Path Cover| = |V| |Maximum Matching| of the new bipartite graph

#### 3 Math

#### 3.1 FFT

```
const int MAXN = 262144;
// (must be 2^k)
// before any usage, run pre_fft() first
typedef long double ld;
typedef complex<ld> cplx;
const ld PI = acosl(-1);
const cplx I(0, 1);
cplx omega[MAXN+1];
void pre_fft(){
  for(int i=0; i<=MAXN; i++)
  omega[i] = exp(i * 2 * PI / MAXN * I);</pre>
// n must be 2^k
void fft(int n, vector<cplx> &a, bool inv=false){
   int basic = MAXN / n;
   int theta = basic;
   for (int m = n; m >= 2; m >>= 1) {
     int mh = m >> 1;
for (int i = 0; i < mh; i++) {</pre>
       cplx w = omega[inv ? MAXN-(i*theta%MAXN)
                              : i*theta%MAXN];
       for (int j = i; j < n; j += m) {
         int k = j + mh;
          cplx x = a[j] - a[k];
         a[j] += a[k];
         a[k] = w * x;
     theta = (theta * 2) % MAXN;
   int i = 0;
   for (int j = 1; j < n - 1; j++) {
  for (int k = n >> 1; k > (i ^= k); k >>= 1);
     if (j < i) swap(a[i], a[j]);</pre>
   if(inv) for (i = 0; i < n; i++) a[i] /= n;
}
```

## 3.2 NTT

```
/* p=a*2^k+1
                                         root
   998244353
                          119
                                  23
   2013265921
                          15
                                  27
                                         31
                                  37
   2061584302081
                          15
   2748779069441
                                  39
   1945555039024054273
                          27
                                  56
                                         5 */
template<ll P,ll root,int MAXK,int MAXN>
struct NTT{
  static ll powi(ll a,ll b){
    ll ret=1;
    for(;b;b>>=1,a=mul(a, a, P)){
      if(b&1) ret=mul(ret, a, P);
    return ret;
  static ll inv(ll a,ll b){
    if(a==1) return 1;
```

```
return (((a-inv(b%a,a))*b+1)/a)%b; // overflow
  11 omega[MAXK+1],inv_omega[MAXK+1];
  NTT(){
    omega[MAXK]=powi(root,(P-1)>>MAXK);
    for(int i=MAXK-1;i>=0;i--)
      omega[i]=mul(omega[i+1], omega[i+1], P);
    for(int i=0;i<=MAXK;i++)</pre>
      inv_omega[i]=inv(omega[i],P);
  void tran(int n,ll a[],bool inv_ntt=false){//n=2^i
    for(int i=1, j=0; i<n; i++){
       for(int k=n>>1;!((j^=k)&k);k>>=1);
      if(i<j) swap(a[i],a[j]);</pre>
    11 *G=(inv_ntt?inv_omega:omega);
    for(int k=2,t=1;k<=n;k<<=1){
  int k2=k>>1;ll dw=G[t++];
       for(int j=0;j<n;j+=k){</pre>
        11 \text{ w=1};
         for(int i=j;i<j+k2;i++){</pre>
           ll x=a[i], y=mul(a[i+k2], w, P);
           a[i]=x+y; if(a[i]>=P) a[i]-=P;
           a[i+k2]=x-y; if(a[i+k2]<0) a[i+k2]+=P;
           w=mul(w, dw, P);
        }
      }
    if(inv_ntt){
      ll inv_n=inv(n,P);
      for(int i=0;i<n;i++) a[i]=mul(a[i], inv_n, P);</pre>
 }
const LL P=2013265921, root=31;
const int MAXN=4194304, MAXK=22; //MAXN=2^k
NTT<P,root,MAXK,MAXN> ntt;
```

#### 3.3 Fast Walsh Transform

```
/* xor convolution:
 * x = (x0,x1) , y = (y0,y1)
* z = (x0y0 + x1y1 , x0y1 + x1y0 )
* x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
* z' = ((x0+x1)(y0+y1), (x0-x1)(y0-y1))
* z = (1/2) * z''
 * or convolution:
 * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
 * and convolution:
 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div
 * ternery xor convolution:
   x = (x0+x1+x2,x0+x1w+x2w^2,x0+x1w^2+x2w)
 * inv = (1/3) * (x0+x1+x2,x0+x1w^2+x2w,x0+x1w+x2w^2)
 * where w^3=1 and w^2=-w-1 */
typedef long long LL;
const int MAXN = (1 << 20) + 10;
const LL MOD = 1e9+7;
inline LL pw( LL x , LL k ) {
  LL res = 1;
  for( LL bs = x ; k ; k >>= 1, bs = (bs * bs)%MOD )
  if( k&1 ) res = ( res * bs ) % MOD;
  return res:
inline LL invf( LL x )
  return pw(x, MOD-2);
inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
  for( int d = 1 ; d < N ; d <<= 1 ) {
    int d2 = d<<1;
}</pre>
     for( int s = 0 ; s < N ; s += d2 )
       for( int i = s , j = s+d ; i < s+d ; i++, j++ ){
 LL ta = x[i] , tb = x[j];
          x[i] = ta+tb;
          x[j] = ta-tb;
if(x[i] >= MOD) x[i] -= MOD;
if(x[j] < 0) x[j] += MOD;
  LL invN = invf(N);
  if( inv )
     for( int i = 0 ; i < N ; i++ ) {</pre>
```

```
x[ i ] *= invN;
x[ i ] %= MOD;
}
```

## 3.4 Poly operator

```
struct PolyOp {
#define FOR(i, c) for (int i = 0; i < (c); ++i)
  NTT<P, root, MAXK, MAXN> ntt;
  static int nxt2k(int x) {
     int i = 1; for (; i < x; i <<= 1); return i;</pre>
  void Mul(int n, LL a[], int m, LL b[], LL c[]) {
   static LL aa[MAXN], bb[MAXN];
     int N = nxt2k(n+m);
    copy(a, a+n, aa); fill(aa+n, aa+N, 0); copy(b, b+m, bb); fill(bb+m, bb+N, 0);
     ntt.tran(N, aa); ntt.tran(N, bb);
    FOR(i, N) c[i] = aa[i] * bb[i] % P;
ntt.tran(N, c, 1);
  void Inv(int n, LL a[], LL b[]) {
    // ab = aa^{-1} = 1 \mod x^{(n/2)}
    // (b - a^-1)^2 = 0 mod x^n
     // bb + a^{-2} - 2 ba^{-1} = 0
     // bba + a^{-1} - 2b = 0
     // a^{-1} = 2b - bba
     static LL tmp[MAXN];
    if (n == 1) {b[0] = ntt.inv(a[0], P); return;}
Inv((n+1)/2, a, b);
     int N = nxt2k(n*2);
    copy(a, a+n, tmp);
fill(tmp+n, tmp+N, 0);
     fill(b+n, b+N, 0);
     ntt.tran(N, tmp); ntt.tran(N, b);
     FOR(i, N) {
       LL t1 = (2 - b[i] * tmp[i]) % P;
       if (t1 < 0) t1 += P
       b[i] = b[i] * t1 % P;
    ntt.tran(N, b, 1);
    fill(b+n, b+N, 0);
  void Div(int n, LL a[], int m, LL b[], LL d[], LL r
       []) {
     // Ra = Rb * Rd mod x^{n-m+1}
     // Rd = Ra * Rb^-1 mod
     static LL aa[MAXN], bb[MAXN], ta[MAXN], tb[MAXN];
     if (n < m) {copy(a, a+n, r); fill(r+n, r+m, 0);</pre>
    // d: n-1 - (m-1) = n-m (n-m+1 terms)
copy(a, a+n, aa); copy(b, b+m, bb);
reverse(aa, aa+n); reverse(bb, bb+m);
Thy(n m+1 bb +b):
    Inv(n-m+1, bb, tb);
Mul(n-m+1, ta, n-m+1, tb, d);
     fill(d+n-m+1, d+n, 0); reverse(d, d+n-m+1);
     // r: m-1 - 1 = m-2 (m-1 terms)
    void dx(int n, LL a[], LL b[]) { REP(i, 1, n-1) b[i
       -1] = i * a[i] % P; }
  void Sx(int n, LL a[], LL b[]) {
    b\lceil 0 \rceil = 0;
     FOR(i, n) b[i+1] = a[i] * ntt.inv(i+1,P) % P;
  void Ln(int n, LL a[], LL b[]) {
   // Integral a' a^-1 dx
     static LL a1[MAXN], a2[MAXN], b1[MAXN];
     int N = nxt2k(n*2)
     dx(n, a, a1); Inv(n, a, a2);
    Mul(n-1, a1, n, a2, b1);
Sx(n+n-1-1, b1, b);
     fill(b+n, b+N, 0);
  void Exp(int n, LL a[], LL b[]) {
    // Newton method to solve g(a(x)) = \ln b(x) - a(x)
     // b' = b - g(b(x)) / g'(b(x))
     // b' = b (1 - lnb + a)
```

```
static LL Inb[MAXN], c[MAXN], tmp[MAXN];
assert(a[0] == 0); // dont know exp(a[0]) mod P
     if (n == 1) {b[0] = 1; return;}
    Exp((n+1)/2, a, b);
fill(b+(n+1)/2, b+n, 0);
     Ln(n, b, lnb);
     fill(c, c+n, 0); c[0] = 1;
     FOR(i, n) {
       c[i] += a[i] - lnb[i];
if (c[i] < 0) c[i] += P
       if (c[i] >= P) c[i] -= P;
    Mul(n, b, n, c, tmp);
     copy(tmp, tmp+n, b);
  bool Sqrt(int n, LL a[], LL b[]){
    // Square root of a : b * b = a ( mod x^n )
// bb = a mod x^n
     // ( bb - a )^2 = 0 mod x^n
     // ( bb + a )^2 = 4 bba
    // ( ( bb + a ) / 2b )^2 = a // sqrt(a) = <math>b / 2 + a / 2b
     static LL c[MAXN];
     int ind=0,x,y,p=1
    while(a[ind]==0) ind++;
     for(int i=0;i<n;i++) a[i]=a[i+ind];</pre>
     if((ind&1)||!solve(a[0],mod,x,y)) // discrete sqrt
        return 0;
    b[0]=min(x,y);
    while(p<n) p<<=1;
for(int t=2;t<=p;t<<=1){</pre>
       Inv(t,b,c); Mul(t,a,t,c,c);
       for(int i=0;i<t;i++)</pre>
          b[i]=(b[i]+c[i])*inv(2)%mod;
     if(ind){
       for(int i=p-1;i>=ind/2;i--) b[i]=b[i-ind/2];
       for(int i=0;i<ind/2;i++) b[i]=0;</pre>
       for(int i=p-1;i>=ind;i--) \bar{a}[\bar{i}]=\bar{a}[i-ind];
       for(int i=0;i<ind;i++) a[i]=0;</pre>
    }
} polyop;
```

## 3.5 Linear Recurrence

```
// Usage: linearRec({0, 1}, {1, 1}, k) //k'th fib
typedef vector<ll> Poly;
ll linearRec(Poly& S, Poly& tr, ll k) {
   int n = tr.size();
  auto combine = [&](Poly& a, Poly& b) {
  Poly res(n * 2 + 1);
  rep(i,0,n+1) rep(j,0,n+1)
     res[i+j]=(res[i+j] + a[i]*b[j])%mod;
for(int i = 2*n; i > n; --i) rep(j,0,n)
res[i-1-j]=(res[i-1-j] + res[i]*tr[j])%mod;
      res.resize(n + 1);
   return res;
}; // a * b mod (x^n-tr)
  Poly pol(n + 1), e(pol);
  pol[0] = e[1] = 1;
   for (++k; k; k /= 2) {
   if (k % 2) pol = combine(pol, e);
      e = combine(e, e);
  11 \text{ res} = 0;
   rep(i,0,n) res=(res + pol[i+1]*S[i])%mod;
   return res:
}
```

## 3.6 BerlekampMassey

```
// find shortest linear recurrence relation O(n^2)
// example: BM({1,1,2,3,5,8,13,21})
// 2*len terms for uniqueness
inline vector<ll> BM(const vector<ll> &x) {
  vector<ll> ls, cur;
int lf; ll ld;
for(int i=0;i<x.size();++i) {</pre>
    ll t=0;
    for(int j=0;j<cur.size();++j)</pre>
       t=(t+x[i-j-1]*cur[j])%mod;
```

```
if((t-x[i])%mod==0) continue;
  if(!cur.size()) {
    cur.resize(i+1);lf=i;ld=(t-x[i])%mod;continue;
  il k=-(x[i]-t)*inv(ld, mod)%mod;
vector<ll> c(i-lf-1); c.push_back(k);
  for(auto j:ls) c.push_back(-j*k%mod);
  if(c.size()<cur.size()) c.resize(cur.size());</pre>
  for(int j=0;j<cur.size();++j)</pre>
    c[j]=(c[j]+cur[j])%mod;
  if(i-lf+(int)ls.size()>=(int)cur.size())
    ls=cur,lf=i,ld=(t-x[i])%mod;
  cur=move(c);
for(auto& xx:cur) xx=(xx%mod+mod)%mod;
return cur;
    Miller Rabin
```

#### 3.7

```
// n < 4,759,123,141
// n < 1,122,004,669,633
                             4:
                                  2, 13, 23, 1662803
// n < 3,474,749,660,383
                                   6 : pirmes <= 13
// n < 2^64
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
bool witness(LL a, LL n, LL u, int t){
  if(!a) return 0;
  LL x=mypow(a,u,n);
  for(int i=0;i<t;i++) {</pre>
    LL nx=mul(x,x,n);
    if(nx==1&&x!=1&&x!=n-1) return 1;
  }
  return x!=1;
bool miller_rabin(LL n,int s=100) {
  // iterate s times of witness on n
  // return 1 if prime, 0 otherwise
  if(n<2) return 0;</pre>
  if(!(n\&1)) return n == 2;
  LL u=n-1; int t=0;
  while(!(u&1)) u>>=1, t++;
  while(s--){
    LL a=randll()%(n-1)+1;
    if(witness(a,n,u,t)) return 0;
  return 1;
}
```

#### 3.8 Simplex

```
/*target:
 \max \sum_{j=1}^n A_{0,j}*x_j
condition:
  \sum_{j=1}^n A_{i,j}*x_j \le A_{i,0} i=1~m
  x_j >= 0 | j=1\sim n
VDB = vector<double>*/
template<class VDB>
VDB simplex(int m,int n,vector<VDB> a){
  vector<int> left(m+1), up(n+1);
  iota(left.begin(), left.end(), n);
  iota(up.begin(), up.end(), 0);
auto pivot = [&](int x, int y){
     swap(left[x], up[y]);
     auto k = a[x][y]; a[x][y] = 1;
     vector<int> pos;
     for(int j = 0; j <= n; ++j){
       a[x][j] /= k;
       if(a[x][j] != 0) pos.push_back(j);
     for(int i = 0; i <= m; ++i){
       if(a[i][y]==0 || i == x) continue;
k = a[i][y], a[i][y] = 0;
       for(int j : pos) a[i][j] -= k*a[x][j];
  for(int x,y;;){
  for(int i=x=1; i <= m; ++i)
    if(a[i][0]<a[x][0]) x = i;
     if(a[x][0] >= 0) break;
     for(int j=y=1; j <= n; ++j)
  if(a[x][j] < a[x][y]) y = j;</pre>
```

|LL solve(LL x1, LL m1, LL x2, LL m2) {

```
if(a[x][y]>=0) return VDB();//infeasible
                                                                      LL g = __gcd(m1, m2);
if((x2 - x1) % g) return -1;// no sol
    pivot(x, y);
                                                                      m1 /= g; m2 /= g;
                                                                      pair<LL,LL> p = gcd(m1, m2);
LL lcm = m1 * m2 * g;
LL res = p.first * (x2 - x1) * m1 + x1;
  for(int x,y;;){
    for(int j=y=1; j <= n; ++j)
if(a[0][j] > a[0][y]) y = j;
    if(a[0][y]<=0) break;
                                                                      return (res % lcm + lcm) % lcm;
    x = -1
    3.11 Pollard Rho
                                                                    // does not work when n is prime
    if(x == -1) return VDB();//unbounded
                                                                    LL f(LL x, LL mod){ return add(mul(x,x,mod),1,mod); }
    pivot(x, y);
                                                                    LL pollard_rho(LL n) {
                                                                      if(!(n&1)) return 2;
  VDB ans(n + 1);
                                                                      while(true){
  for(int i = 1; i <= m; ++i)
                                                                        LL y=2, x=rand()\%(n-1)+1, res=1;
    if(left[i] \le n) ans[left[i]] = a[i][0];
                                                                        for(int sz=2; res==1; sz*=2) {
  ans[0] = -a[0][0];
                                                                           for(int i=0; i<sz && res<=1; i++) {</pre>
  return ans;
                                                                             x = f(x, n)
                                                                             res = \_gcd(abs(x-y), n);
3.9 Faulhaber
                                                                          y = x;
 * faulhaber's formula -
                                                                        if (res!=0 && res!=n) return res;
 * cal power sum formula of all p=1\simk in O(k^2) */
#define MAXK 2500
                                                                   }
const int mod = 1000000007;
int b[MAXK]; // bernoulli number
                                                                    3.12 ax+by=gcd
int inv[MAXK+1]; // inverse
int cm[MAXK+1][MAXK+1]; // combinactories
int co[MAXK][MAXK+2]; // coeeficient of x^j when p=i
                                                                   pair<ll,ll> gcd(ll a, ll b){
                                                                      if(b == \emptyset) return {1, \emptyset};
inline int getinv(int x) {
                                                                      pair<ll,ll>q=gcd(b,a\%b);
  int a=x,b=mod,a0=1,a1=0,b0=0,b1=1;
                                                                      return {q.second, q.first - q.second * (a / b)};
  while(b) {
                                                                   }
    int q,t;
                                                                    3.13 Discrete sqrt
    q=a/b; t=b; b=a-b*q; a=t;
    t=b0; b0=a0-b0*q; a0=t;
                                                                   void calcH(int &t, int &h, const int p) {
    t=b1; b1=a1-b1*q; a1=t;
                                                                      int tmp=p-1; for(t=0;(tmp&1)==0;tmp/=2) t++; h=tmp;
  return a0<0?a0+mod:a0;</pre>
                                                                    // solve equation x^2 \mod p = a where p is a prime
                                                                   bool solve(int a, int p, int &x, int &y) {
inline void pre() {
  /* combinational */
                                                                      if(p == 2) { x = y = 1; return true; }
                                                                      int p2 = p / 2, tmp = mypow(a, p2, p);
if (tmp == p - 1) return false;
  for(int i=0;i<=MAXK;i++) {</pre>
    cm[i][0]=cm[i][i]=1;
                                                                      if ((p + 1) \% 4 == 0) {
    for(int j=1;j<i;j++)
  cm[i][j]=add(cm[i-1][j-1],cm[i-1][j]);</pre>
                                                                        x=mypow(a,(p+1)/4,p); y=p-x; return true;
                                                                      } else {
                                                                        int t, h, b, pb; calcH(t, h, p);
  /* inverse */
                                                                        if (t >= 2) {
  for(int i=1;i<=MAXK;i++) inv[i]=getinv(i);</pre>
                                                                           do \{b = rand() \% (p - 2) + 2;
  /* bernoulli */
                                                                           } while (mypow(b, p / 2, p) != p - 1);
  b[0]=1; b[1]=getinv(2); // with b[1] = 1/2 for(int i=2;i<MAXK;i++) {
                                                                        pb = mypow(b, h, p);

} int s = mypow(a, h / 2, p);

for (int step = 2; step <= t; step++) {
    if(i&1) { b[i]=0; continue; }
    b[i]=1;
                                                                           int ss = (((LL)(s * s) % p) * a) % p;
    for(int j=0;j<i;j++)</pre>
                                                                           for(int i=0;i<t-step;i++) ss=mul(ss,ss,p);</pre>
       b[i]=sub(b[i], mul(cm[i][j], mul(b[j], inv[i-j+1])
                                                                          if (ss + 1 == p) s = (s * pb) % p;
pb = ((LL)pb * pb) % p;
           ));
                                                                        x = ((LL)s * a) % p; y = p - x;
  /* faulhaber */
                                                                      } return true;
  // sigma_x=1^n {x^p} = // 1/(p+1) * sigma_j=0^p {C(p+1,j)*Bj*n^(p-j+1)}
                                                                   }
  for(int i=1;i<MAXK;i++) {</pre>
                                                                    3.14 Romberg
    co[i][0]=0;
    for(int j=0;j<=i;j++)
  co[i][i-j+1]=mul(inv[i+1], mul(cm[i+1][j], b[j]))</pre>
                                                                    // Estimates the definite integral of
                                                                    // \cdot int_a^b f(x) dx
                                                                    template<class T>
                                                                    double romberg( T& f, double a, double b, double eps=1e
                                                                         -8){
/* sample usage: return f(n,p) = sigma_x=1\sim (x^p) */
                                                                      vector<double>t; double h=b-a,last,curr; int k=1,i=1;
t.push_back(h*(f(a)+f(b))/2);
inline int solve(int n,int p) {
  int sol=0,m=n;
                                                                      do{ last=t.back(); curr=0; double x=a+h/2;
  for(int i=1;i<=p+1;i++) {</pre>
                                                                        for(int j=0;j<k;j++) curr+=f(x), x+=h;
curr=(t[0] + h*curr)/2; double k1=4.0/3.0,k2</pre>
    sol=add(sol,mul(co[p][i],m));
    m = mul(m, n);
                                                                             =1.0/3.0;
                                                                        for(int j=0;j<i;j++){ double temp=k1*curr-k2*t[j];</pre>
                                                                        t[j]=curr; curr=temp; k2/=4*k1-k2; k1=k2+1; 
t.push_back(curr); k*=2; h/=2; i++;
  return sol;
                                                                      }while( fabs(last-curr) > eps);
3.10 Chinese Remainder
                                                                      return t.back();
```

#### 3.15 Prefix Inverse

```
void solve( int m ){
  inv[ 1 ] = 1;
  for( int i = 2 ; i < m ; i ++ )
     inv[ i ] = ((LL)(m - m / i) * inv[m % i]) % m;
}</pre>
```

## 3.16 Roots of Polynomial

```
const double eps = 1e-12;
const double inf = 1e+12;
double a[ 10 ], x[ 10 ];
int sign( double x ){return (x < -eps)?(-1):(x>eps);}
double f(double a[], int n, double x){
  double tmp=1,sum=0;
for(int i=0;i<=n;i++)</pre>
  { sum=sum+a[i]*tmp; tmp=tmp*x; }
  return sum;
double binary(double l,double r,double a[],int n){
  int sl=sign(f(a,n,l)), sr=sign(f(a,n,r));
if(sl==0) return l; if(sr==0) return r;
if(sl*sr>0) return inf;
  while(r-l>eps){
     double mid=(l+r)/2;
     int ss=sign(f(a,n,mid));
     if(ss==0) return mid;
     if(ss*sl>0) l=mid; else r=mid;
  }
  return 1;
void solve(int n,double a[],double x[],int &nx){
  if(n==1){ x[1]=-a[0]/a[1]; nx=1; return; }
  double da[10], dx[10]; int ndx;
for(int i=n;i>=1;i--) da[i-1]=a[i]*i;
  solve(n-1,da,dx,ndx);
  nx=0:
  if(ndx==0){
     double tmp=binary(-inf,inf,a,n);
     if (tmp<inf) x[++nx]=tmp;</pre>
     return;
  double tmp;
tmp=binary(-inf,dx[1],a,n);
  if(tmp<inf) x[++nx]=tmp;</pre>
  for(int i=1;i<=ndx-1;i++){</pre>
     tmp=binary(dx[i],dx[i+1],a,n);
     if(tmp<inf) x[++nx]=tmp;</pre>
  tmp=binary(dx[ndx],inf,a,n);
  if(tmp<inf) x[++nx]=tmp;</pre>
int main() {
  scanf("%d",&n);
  for(int i=n;i>=0;i--) scanf("%lf",&a[i]);
  int nx:
  solve(n,a,x,nx);
  for(int i=1;i<=nx;i++) printf("%.6f\n",x[i]);</pre>
```

## 3.17 Primes and $\mu$ function

```
/* 12721, 13331, 14341, 75577, 123457, 222557, 556679
* 999983, 1097774749, 1076767633, 100102021, 999997771
* 1001010013, 1000512343, 987654361, 999991231
* 999888733, 98789101, 987777733, 999991921, 1010101333
* 1010102101, 1000000000039, 100000000000037
* 2305843009213693951, 4611686018427387847
* 9223372036854775783, 18446744073709551557 */
int mu[ N ] , p_tbl[ N ]; // multiplicative function f
vector-int> primes;
void sieve() {
   mu[ 1 ] = p_tbl[ 1 ] = 1;
   for( int i = 2 ; i < N ; i ++ ){
      if( !p_tbl[ i ] ){
        p_tbl[ i ] = i;
        primes.push_back( i );
        mu[ i ] = -1; // f(i)=... where i is prime
    }
   for( int p : primes ){</pre>
```

#### 3.18 Result

- Lucas' Theorem : For  $n,m\in\mathbb{Z}^*$  and prime P,  $C(m,n)\mod P=\Pi(C(m_i,n_i))$  where  $m_i$  is the i-th digit of m in base P.
- 1st Stirling Numbers(permutation |P|=n with k cycles): S(n,k)= coefficient of  $x^k$  in  $\Pi_{i=0}^{n-1}(x+i)$  S(n+1,k)=nS(n,k)+S(n,k-1)
- 2nd Stirling Numbers(Partition n elements into k non-empty set):  $S(n,k)=\frac{1}{k!}\sum_{j=0}^k(-1)^{k-j}\binom{k}{j}j^n\\ S(n+1,k)=kS(n,k)+S(n,k-1)$
- Calculate f(x+n) where  $f(x) = \sum\limits_{i=0}^{n-1} a_i x^i$ :  $f(x+n) = \sum\limits_{i=0}^{n-1} a_i (x+n)^i = \sum\limits_{i=0}^{n-1} x^i \cdot \frac{1}{i!} \sum\limits_{j=i}^{n-1} \frac{a_j}{j!} \cdot \frac{n^{j-i}}{(j-i)!}$
- Calculate  $c[i-j]+=a[i]\times b[j]$  for a[n],b[m] 1. a=reverse(a); c=mul(a,b); c=reverse(c[:n]); 2. b=reverse(b); c=mul(a,b); c=rshift(c,m-1);
- Eulerian number(permutation  $1\sim n$  with m a[i]>a[i-1]):  $A(n,m)=\sum_{i=0}^m (-1)^i {n+1\choose i}(m+1-i)^n$  A(n,m)=(n-m)A(n-1,m-1)+(m+1)A(n-1,m)
- Derangement:  $D(n) = (n-1)(D(n-1) + D(n-2)) = nD(n-1) + (-1)^n$
- Pick's Theorem : A=i+b/2-1
- Kirchhoff's theorem : number of spanning tree of undirected graph: degree matrix  $D_{ii} = deg(i)$  ,  $D_{ij} = 0$  adjacency matrix  $G_{ij} = \# \ of \ (i,j) \in E$  ,  $G_{ii} = 0$ , let A = D G, delete any one row, one column, and cal det(A') number of spanning tree of directed graph: in-degree matrix  $D_{ii}^{in} = indeg(i)$  ,  $D_{ij}^{in} = 0$  out-degree matrix  $D_{ii}^{out} = outdeg(i)$  ,  $D_{ij}^{out} = 0$  let  $L^{in} = D^{in} G$  ,  $L^{out} = D^{out} G$  , delete the i-th row and column  $det(L_i^{in})$  and  $det(L_i^{out})$  is the number of spanning tree from/to root i
- Burnside Lemma:  $|X/G| = \frac{1}{|G|} \sum\limits_{g \in G} |X^g|$
- Polya theorem:  $|Y^x/G|=\frac{1}{|G|}\sum_{g\in G}m^{c(g)}$  m=|Y| : num of colors, c(g) : num of cycle
- Anti SG (the person who has no strategy wins) : first player wins iff either 1. SG value of ALL subgame  $\leq 1$  and SG value of the game = 0 2. SG value of some subgame > 1 and SG value of the game  $\neq 0$
- Möbius inversion formula :  $g(n) = \sum_{d \mid n} f(d) \text{ for every integer } n \geq 1 \text{ , then}$   $f(n) = \sum_{d \mid n} \mu(d) g(\frac{n}{d}) = \sum_{d \mid n} \mu(\frac{n}{d}) g(d) \text{ for every integer } n \geq 1$  Dirichlet convolution :  $f * g = g * f = \sum_{d \mid n} f(d) g(\frac{n}{d}) = \sum_{d \mid n} f(\frac{n}{d}) g(d)$   $g = f * 1 \Leftrightarrow f = g * \mu, \; \epsilon = \mu * 1, \; Id = \phi * 1, \; d = 1 * 1, \; \sigma = Id * 1 = \phi * d, \\ \sigma_k = Id_k * 1 \text{ where } \epsilon(n) = [n-1], \; 1(n) = 1, \; Id(n) = n, \; Id_k(n) = n^k,$

d(n) = #(divisor),  $\sigma(n) = \sum divisor$ ,  $\sigma_k(n) = \sum divisor^k$ 

```
• Find a Primitive Root of n: n has primitive roots iff n=2,4,p^k,2p^k where p is an odd prime. 1. Find \phi(n) and all prime factors of \phi(n), says P=\{p_1,...,p_m\} 2. \forall g \in [2,n), if g^{\frac{\phi(n)}{p_i}} \neq 1, \forall p_i \in P, then g is a primitive root. 3. Since the smallest one isn't too big, the algorithm runs fast. 4. n has exactly \phi(\phi(n)) primitive roots.

• Sum of Two Squares Thm (Legendre): For a given positive integer N, let D1=(\#\ of\ d\in N\ dividing\ N\ that\ d=1\ (mod\ 4)) D3=(\#\ of\ d\in N\ dividing\ N\ that\ d=3\ (mod\ 4)) then N can be written as a sum of two squares in exactly R(N)=4(D1-D3) ways.

• Difference of R(N)=4(D1-D3) ways.

• Difference of R(N)=4(D1-D3) hm: let R(N)=4(D1-D3) hm: let R(N)=4(D1-D3) hm: R(N)
```

# 4 Geometry

## 4.1 Intersection of 2 lines

```
Pt LLIntersect(Line a, Line b) {
  Pt p1 = a.s, p2 = a.e, q1 = b.s, q2 = b.e;
  ld f1 = (p2-p1)^(q1-p1),f2 = (p2-p1)^(p1-q2),f;
  if(dcmp(f=f1+f2) == 0)
    return dcmp(f1)?Pt(NAN,NAN):Pt(INFINITY,INFINITY);
  return q1*(f2/f) + q2*(f1/f);
}
```

## 4.2 halfPlaneIntersection

```
^{\prime} for point or line solution, change > to >=
bool onleft(Line L, Pt p) {
  return dcmp(L.v^{p-L.s}) > 0;
// assume that Lines intersect
vector<Pt> HPI(vector<Line>& L) {
  sort(L.begin(), L.end());
int n = L.size(), fir, las;
  Pt *p = new Pt[n];
  Line *q = new Line[n];
  q[fir=las=0] = L[0];
  for(int i = 1; i < n; i++) {
    while(fir < las && !onleft(L[i], p[las-1])) las--;</pre>
    while(fir < las && !onleft(L[i], p[fir])) fir++;</pre>
    q[++las] = L[i];
    if(dcmp(q[las].v^q[las-1].v) == 0) {
      las-
      if(onleft(q[las], L[i].s)) q[las] = L[i];
    if(fir < las) p[las-1] = LLIntersect(q[las-1], q[</pre>
        las]);
  while(fir < las && !onleft(q[fir], p[las-1])) las--;</pre>
  if(las-fir <= 1) return {};</pre>
  p[las] = LLIntersect(q[las], q[fir]);
  int m = 0;
  vector<Pt> ans(las-fir+1);
  for(int i = fir ; i <= las ; i++) ans[m++] = p[i];</pre>
  return ans;
```

#### 4.3 Intersection of 2 segments

```
if(isfinite(p.x) && onseg(p, a) && onseg(p, b))
    return p; //not parallel
  return {NAN,NAN};
}
```

#### 4.4 Banana

#### 4.5 Intersection of circle and line

```
vector<Pt> CLInter(const Line &a,const Circle &c){
   Pt p=a.s+(c.o-a.s)*a.v/norm2(a.v)*a.v;
   ld d=c.r*c.r-norm2(c.o-p);
   if(d<-eps) return {};
   if(d<eps) return {p};
   Pt v=a.v/norm(a.v)*sqrt(d);
   return {p+v,p-v};
}</pre>
```

## 4.6 Intersection of polygon and circle

```
ld PCIntersect(vector<Pt> v, Circle cir) {
  for(int i = 0 ; i < (int)v.size() ; ++i) v[i] = v[i]</pre>
        - cir.o;
  ld ans = 0, r = cir.r;
  int n = v.size();
  for(int i = 0; i < n; ++i) {
Pt pa = v[i], pb = v[(i+1)%n];
     if(norm(pa) < norm(pb)) swap(pa, pb);</pre>
     if(dcmp(norm(pb)) == 0) continue;
     ld s, h, theta;
     ld a = norm(pb), b = norm(pa), c = norm(pb-pa);
    ld cosB = (pb*(pb-pa))/a/c, B = acos(cosB);
if(cosB > 1) B = 0;
     else if(cosB < -1) B = PI;</pre>
     ld cosC = (pa*pb)/a/b, C = acos(cosC);
     if(cosC > 1) C = 0;
     else if(cosC < -1) C = PI;</pre>
     if(a > r) {
       s = (C/2)*r*r
       h = a*b*sin(c)/c;
if(h < r && B < PI/2) s -= (acos(h/r)*r*r - h*
            sqrt(r*r-h*h));
     else if(b > r) {
  theta = PI - B - asin(sin(B)/r*a);
       s = 0.5*a*r*sin(theta) + (C-theta)/2*r*r;
     else s = 0.5*sin(C)*a*b;
    ans += abs(s)*dcmp(v[i]^v[(i+1)%n]);
  return abs(ans);
```

# 4.7 Intersection of 2 circles

## 4.8 Circle cover

```
#define N 1021
struct CircleCover{
  int C; Circle c[N];
  bool g[N][N], overlap[N][N];
  // Area[i] : area covered by at least i circles
  ld Area[N];
  void init(int _C){ C = _C; }
  bool CCinter(Circle& a, Circle& b, Pt& p1, Pt& p2){
```

```
Pt o1 = a.o, o2 = b.o; ld r1 = a.r, r2 = b.r; if(norm(o1 - o2) > r1 + r2) return 0; if(norm(o1 - o2) < max(r1, r2) - min(r1, r2))
           return 0;
     1d d2 = (o1 - o2) * (o1 - o2);
     ld d = sqrt(d2);
     if(d > r1 + r2) return 0;
Pt u=(o1+o2)*0.5 + (o1-o2)*((r2*r2-r1*r1)/(2*d2));
     1d A=sqrt((r1+r2+d)*(r1-r2+d)*(r1+r2-d)*(-r1+r2+d))
     Pt v=Pt(o1.y-o2.y, -o1.x + o2.x) * A / (2*d2);
     p1 = u + v; p2 = u - v;
     return 1;
   struct Teve {
     Pt p; ld ang; int add;
     Teve() {}
     Teve(Pt _a, ld _b, int _c):p(_a), ang(_b), add(_c)
     bool operator<(const Teve &a)const
     {return ang < a.ang;}
  }eve[N * 2];
  // strict: x = 0, otherwise x = -1
bool disjuct(Circle& a, Circle &b, int x)
   {return sign(norm(a.o - b.o) - a.r - b.r) > x;}
   bool contain(Circle& a, Circle &b, int x)
   {return sign(a.r - b.r - norm(a.o - b.o)) > x;}
   bool contain(int i, int j){
     /* c[j] is non-strictly_in c[i]. */
     return (sign(c[i].r - c[j].r) > 0 ||
(sign(c[i].r - c[j].r) == 0 && i < j)) &&
                     contain(c[i], c[j], -1);
   void solve(){
     for(int i = 0; i <= C + 1; i++) Area[i] = 0;
     for(int i = 0; i < C; i++)
  for(int j = 0; j < C; j++)
   g[i][j] = !(overlap[i][j] || overlap[j][i] ||</pre>
                         disjuct(c[i], c[j], -1));
     for(int i = 0; i < C; i++){
        int E = 0, cnt = 1;
        for(int j = 0; j < C; j++)
  if(j != i && overlap[j][i])</pre>
             cnt++;
        for(int j = 0; j < C; j++)</pre>
          if(i != j && g[i][j]){
Pt aa, bb;
             CCinter(c[i], c[j], aa, bb);
             ld A=atan2(aa.y - c[i].o.y, aa.x - c[i].o.x);
ld B=atan2(bb.y - c[i].o.y, bb.x - c[i].o.x);
             eve[E++] = Teve(bb, B, 1);
             eve[E++] = Teve(aa, A, -1);
             if(B > A) cnt++;
        if(E == 0) Area[cnt] += pi * c[i].r * c[i].r;
        else{
          sort(eve , eve_+ E);
          eve[E] = eve[0];
          for(int j = 0; j < E; j++){
  cnt += eve[j].add;</pre>
             Area[cnt] += (eve[j].p \wedge eve[j + 1].p) * .5;
             ld theta = eve[j + 1].ang - eve[j].ang;
if (theta < 0) theta += 2. * pi;</pre>
             Area[cnt] +=
               (theta - sin(theta)) * c[i].r*c[i].r * .5;
```

#### 4.9 Li Chao Segment Tree

```
struct LiChao_min{
    struct line{
        LL m, c;
        line(LL _m=0, LL _c=0) { m = _m; c = _c; }
        LL eval(LL x) { return m * x + c; }
    };
    struct node{
        node *l, *r; line f;
        node(line v) { f = v; l = r = NULL; }
    };
```

```
typedef node* pnode;
pnode root; int sz;
#define mid ((l+r)>>1)
   void insert(line &v, int l, int r, pnode &nd){
     if(!nd) { nd = new node(v); return; }
     LL trl = nd->f.eval(l), trr = nd->f.eval(r);
     LL vl = v.eval(l), vr = v.eval(r);
     if(trl <= vl && trr <= vr) return;</pre>
     if(trl > vl && trr > vr) { nd->f = v; return; }
     if(trl > vl) swap(nd->f, v);
     if(nd->f.eval(mid) < v.eval(mid)) insert(v, mid +</pre>
          1, r, nd->r);
     else swap(nd->f, v), insert(v, l, mid, nd->l);
   LL query(int x, int 1, int r, pnode &nd){
     if(!nd) return LLONG_MAX;
     if(l == r) return nd->f.eval(x);
     if(mid  >=  x ) return min(nd- >f.eval(x), query(x, l,
          mid, nd->1));
     return min(nd->f.eval(x), query(x, mid + 1, r, nd->
          r));
  /* -sz <= query_x <= sz */
void init(int _sz){ sz = _sz + 1; root = NULL; }
void add_line(LL m, LL c){ line v(m, c); insert(v, -</pre>
        sz, sz, root); }
  LL query(LL x) { return query(x, -sz, sz, root); }
};
```

```
4.10 Convex Hull trick
/* Given a convexhull, answer querys in O(\lg N) CH should not contain identical points, the area should be > 0, min pair(x, y) should be listed first */ double det( const Pt& p1 , const Pt& p2 )
{ return p1.x * p2.y - p1.y * p2.x; }
struct Conv{
   int n;
   vector<Pt> a;
   vector<Pt> upper, lower;
   Conv(vector < Pt > \_a) : a(\_a){}
     n = a.size();
     int ptr = 0;
     for(int i=1; i<n; ++i) if (a[ptr] < a[i]) ptr = i;
for(int i=0; i<=ptr; ++i) lower.push_back(a[i]);</pre>
     for(int i=ptr; i<n; '++i) upper.push_back(a[i]);
upper.push_back(a[0]);</pre>
   int sign( LL x ){ // fixed when changed to double
  return x < 0 ? -1 : x > 0; }
   pair<LL,int> get_tang(vector<Pt> &conv, Pt vec){
     int l = 0, r = (int)conv.size() - 2;
     while(l + 1 < r){
        int mid = (l + r) / 2;
        if(sign(det(conv[mid+1]-conv[mid],vec))>0)r=mid;
        else l = mid;
     void upd_tang(const Pt &p, int id, int &i0, int &i1){
  if(det(a[i0] - p, a[id] - p) > 0) i0 = id;
  if(det(a[i1] - p, a[id] - p) < 0) i1 = id;</pre>
   void bi_search(int l, int r, Pt p, int &i0, int &i1){
     if(l == r) return;
     upd_tang(p, 1 % n, i0, i1);
     int sl=sign(det(a[l % n] - p, a[(l + 1) % n] - p));
     while(l + 1 < r) {
  int mid = (l + r) / 2;</pre>
        int smid=sign(det(a[mid%n]-p, a[(mid+1)%n]-p));
        if (smid == sl) l = mid;
        else r = mid;
     upd_tang(p, r % n, i0, i1);
   int bi_search(Pt u, Pt v, int l, int r) .
     int sl = sign(det(v - u, a[1 % n] - u));
     while(l + 1 < r) 
        int mid = (l + r) / 2;
int smid = sign(det(v - u, a[mid % n] - u));
        if (smid == sl) l = mid;
```

```
else r = mid:
    return 1 % n;
  // 1. whether a given point is inside the CH
  bool contain(Pt p) {
    if (p.x < lower[0].x | l p.x > lower.back().x)
         return 0;
    int id = lower_bound(lower.begin(), lower.end(), Pt
         (p.x, -INF)) - lower.begin();
    if (lower[id].x == p.x) {
      if (lower[id].y > p.y) return 0;
    }else if(det(lower[id-1]-p,lower[id]-p)<0)return 0;</pre>
    id = lower_bound(upper.begin(), upper.end(), Pt(p.x
          INF), greater<Pt>()) - upper.begin();
    if (upper[id].x == p.x) {
    if (upper[id].y < p.y) return 0;
}else if(det(upper[id-1]-p,upper[id]-p)<0)return 0;</pre>
    return 1;
  }
  // 2. Find 2 tang pts on CH of a given outside point
  // return true with i0, i1 as index of tangent points
  // return false if inside CH
  bool get_tang(Pt p, int &i0, int &i1) {
    if (contain(p)) return false;
    i0 = i1 = 0;
    int id = lower_bound(lower.begin(), lower.end(), p)
          lower.begin();
    bi_search(0, id, p, i0, i1);
bi_search(id, (int)lower.size(), p, i0, i1);
    id = lower_bound(upper.begin(), upper.end(), p,
         greater<Pt>()) - upper.begin();
    bi_search((int)lower.size() - 1, (int)lower.size()
         -1 + id, p, i0, i1);
    bi_search((int)lower.size() - 1 + id, (int)lower.
         size() - 1 + (int)upper.size(), p, i0, i1);
    return true;
  // 3. Find tangent points of a given vector
  // ret the idx of vertex has max cross value with vec
  int get_tang(Pt vec){
    pair<LL, int> ret = get_tang(upper, vec);
    ret.second = (ret.second+(int)lower.size()-1)%n;
    ret = max(ret, get_tang(lower, vec));
    return ret.second;
  // 4. Find intersection point of a given line
  // return 1 and intersection is on edge (i, next(i))
  // return 0 if no strictly intersection
  bool get_intersection(Pt u, Pt v, int &i0, int &i1){
   int p0 = get_tang(u - v), p1 = get_tang(v - u);
   if(sign(det(v-u,a[p0]-u))*sign(det(v-u,a[p1]-u))<0){
      if (p0 > p1) swap(p0, p1);
     i0 = bi_search(u, v, p0, p1);
     i1 = bi\_search(u, v, p1, p0 + n);
     return 1;
   return 0;
};
```

#### 4.11 Tangent line of two circles

```
vector<Line> go(const Circle& c1, const Circle& c2, int
     sign1){
  // sign1 = 1 for outer tang, -1 for inter tang
  vector<Line> ret;
  double d_{sq} = norm2(c1.o - c2.o);
  if(d_sq < eps) return ret;</pre>
  double d = sqrt(d_sq);
 Pt v = (c2.0 - c1.0) / d;
  double c = (c1.r - sign1 * c2.r) / d;
  if(c * c > 1) return ret;
  double h = sqrt(max(0.0, 1.0 - c * c));
  for(int sign2 = 1; sign2 >= -1; sign2 -= 2){
  Pt n = { v.x * c - sign2 * h * v.y ,
              v.y * c + sign2 * h * v.x };
    Pt p1 = c1.o + n * c1.r;
    Pt p2 = c2.0 + n * (c2.r * sign1);
    if(fabs(p1.x - p2.x) < eps and
       fabs(p1.y - p2.y) < eps)
      p2 = p1 + perp(c2.o - c1.o);
```

```
12
    ret.push_back({p1, p2});
  return ret;
}
4.12 Tangent line of point and circle
vector<Line> PCTangent(const Circle& C, const Pt& P) {
  vector<Line> ans;
  Pt u = C.o - P;
  double dist = norm(u);
  if(dist < C.r) return ans;</pre>
  else if(abs(dist) < eps) {</pre>
    ans.push_back({P, P+rotate(u, M_PI/2)});
    return ans;
  else {
    double ang = asin(C.r/dist);
    ans.push_back({P, P+rotate(u, -ang)});
ans.push_back({P, P+rotate(u, +ang)});
    return ans;
4.13 Min distance of two convex
double TwoConvexHullMinDis(Point P[], Point Q[], int n,
      int m) {
  int YMinP=0, YMaxQ=0; double tmp, ans=1e9;
for(int i=0;i<n;++i) if(P[i].y<P[YMinP].y) YMinP=i;
for(int i=0;i<m;++i) if(Q[i].y>Q[YMaxQ].y) YMaxQ=i;
  P[n]=P[0]; Q[m]=Q[0];
  for (int i=0;i<n;++i) {</pre>
    while(tmp=((Q[YMaxQ+1]-P[YMinP+1])^(P[YMinP]-P[
         YMinP+1]))>((Q[YMaxQ]-P[YMinP+1])^(P[YMinP]-P[
         YMinP+1]))) YMaxQ=(YMaxQ+1)%m;
    if(tmp<0)ans=min(ans,PtToSegDis(P[YMinP],P[YMinP</pre>
         +1],Q[YMaxQ]));
    else ans=min(ans,TwoSegMinDis(P[YMinP],P[YMinP+1],Q
         [YMaxQ],Q[YMaxQ+1]));
    YMinP=(YMinP+1)%n;
  return ans;
4.14 Poly Union
struct PY{
  int n; Pt pt[5]; double area;
  Pt& operator[](const int x){ return pt[x]; }
  void init(){ //n,pt[0~n-1] must be filled
    area=pt[n-1]^pt[0];
     for(int i=0;i<n-1;i++) area+=pt[i]^pt[i+1];
     if((area/=2)<0)reverse(pt,pt+n),area=-area;</pre>
PÝ py[500];
pair<double,int> c[5000];
inline double segP(Pt &p,Pt &p1,Pt &p2){
  if(dcmp(p1.x-p2.x)==0) return (p.y-p1.y)/(p2.y-p1.y);
  return (p.x-p1.x)/(p2.x-p1.x);
double polyUnion(int n){ //py[0~n-1] must be filled
  int i,j,ii,jj,ta,tb,r,d;
  double z,w,s,sum,tc,td
  for(i=0;i<n;i++) py[i][py[i].n]=py[i][0];</pre>
  sum=0;
  for(i=0;i< n;i++)\{
    for(ii=0;ii<py[i].n;ii++){</pre>
       c[r++]=make_pair(0.0,0);
       c[r++]=make_pair(1.0,0);
       for(j=0;j<n;j++){</pre>
         if(i==j) continue;
         for(jj=0;jj<py[j].n;jj++){</pre>
           ta=dcmp(tri(py[i][ii],py[i][ii+1],py[j][jj]))
           tb=dcmp(tri(py[i][ii],py[i][ii+1],py[j][jj
               +1])):
```

if(ta==0 && tb==0){

i][ii])>0 && j<i){

if((py[j][jj+1]-py[j][jj])\*(py[i][ii+1]-py[

```
National Taiwan University CRyptoGRapheR
                c[r++]=make_pair(segP(py[j][jj],py[i][ii
                     ],py[i][ii+1]),<u>1</u>)
                c[r++j=make_pair(segP(py[j][jj+1],py[i][
                     ii],py[i][ii+1]),-1);
           }else if(ta>=0 && tb<0){
              tc=tri(py[j][jj],py[j][jj+1],py[i][ii]);
td=tri(py[j][jj],py[j][jj+1],py[i][ii+1]);
              c[r++]=make_pair(tc/(tc-td),1);
           }else if(ta<0 && tb>=0){
              tc=tri(py[j][jj],py[j][jj+1],py[i][ii]);
td=tri(py[j][jj],py[j][jj+1],py[i][ii+1]);
              c[r++]=make_pair(tc/(tc-td),-1);
         }
       sort(c,c+r);
       z=min(max(c[0].first,0.0),1.0);
       d=c[0].second; s=0;
       for(j=1; j<r; j++){</pre>
         w=min(max(c[j].first,0.0),1.0);
         if(!d) s+=w-z
         d+=c[j].second; z=w;
       sum+=(py[i][ii]^py[i][ii+1])*s;
  return sum/2;
}
4.15 Lower Concave Hull
const ll is_query = -(1LL<<62);</pre>
struct Line {
  11 m, b;
  mutable function<const Line*()> succ;
  bool operator<(const Line& rhs) const {</pre>
    if (rhs.b != is_query) return m < rhs.m;</pre>
    const Line* s = succ();
```

```
return s ? b - s->b < (s->m - m) * rhs.m : 0;
}; // maintain upper hull for maximum
struct HullDynamic : public multiset<Line> {
  bool bad(iterator y) {
    auto z = next(y);
    if (y == begin()) {
      if (z == end()) return 0;
      return y->m == z->m && y->b <= z->b;
    auto x = prev(y);
    if(z==end())return y->m==x->m&y->b<=x->b;
    return (x->b-y->b)*(z->m-y->m)>=
(y->b-z->b)*(y->m-x->m);
  void insert_line(ll m, ll b) {
  auto y = insert({m, b});
    y->succ = [=]{return next(y)==end()?0:&*next(y);};
    if(bad(y)) {erase(y); return; }
    while(next(y)!=end()&&bad(next(y)))erase(next(y));
    while(y!=begin()&&bad(prev(y)))erase(prev(y));
  ll eval(ll x) {
    auto l = *lower_bound((Line) {x, is_query});
     return l.m * x + l.b;
};
```

## 4.16 Delaunay Triangulation

```
/* Delaunay Triangulation:
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle.

find : return a triangle contain given point
add_point : add a point into triangulation

A Triangle is in triangulation iff. its has_chd is 0.
Region of triangle u: iterate each u.edge[i].tri,
each points are u.p[(i+1)%3], u.p[(i+2)%3] */
const int N = 100000 + 5;
```

```
const type inf = 2e3;
type eps = 1e-6; // 0 when integer
type sqr(type x) { return x*x; }
  return p4 is in circumcircle of tri(p1,p2,p3)
bool in_cc(const Pt& p1, const Pt& p2, const Pt& p3,
     const Pt& p4){
  type u11 = p\dot{1}.\dot{X} - p4.X; type u12 = p1.Y - p4.Y; type u21 = p2.X - p4.X; type u22 = p2.Y - p4.Y;
  type u31 = p3.X - p4.X; type u32 = p3.Y - p4.Y;
  type u13 = sqr(p1.X)-sqr(p4.X)+sqr(p1.Y)-sqr(p4.Y);
  type u23 = sqr(p2.X)-sqr(p4.X)+sqr(p2.Y)-sqr(p4.Y);
  type u33 = sqr(p3.X)-sqr(p4.X)+sqr(p3.Y)-sqr(p4.Y);
  type det = -u13*u22*u31 + u12*u23*u31 + u13*u21*u32
               -u11*u23*u32 - u12*u21*u33 + u11*u22*u33;
  return det > eps;
type side(const Pt& a, const Pt& b, const Pt& p)
{ return (b - a) ^ (p - a); }
typedef int SdRef;
struct Tri;
typedef Tri* TriRef;
struct Edge {
  TriRef tri; SdRef side;
Edge():tri(0), side(0){}
  Edge(TriRef _tri, SdRef _side):tri(_tri), side(_side)
       {}
};
struct Tri {
  Pt p[3];
  Edge edge[3]
  TriRef chd[3];
  Tri() {}
  Tri(const Pt& p0, const Pt& p1, const Pt& p2) {
  p[0] = p0; p[1] = p1; p[2] = p2;
     chd[0] = chd[1] = chd[2] = 0;
  bool has_chd() const { return chd[0] != 0; }
  int num_chd() const {
    return chd[0] == 0? 0
          : chd[1] == 0 ? 1
          : chd[2] == 0 ? 2 : 3;
  bool contains(Pt const& q) const {
    for( int i = 0; i < 3; i ++ )
       if( side(p[i], p[(i + 1) % 3] , q) < -eps )
         return false;
    return true;
} pool[ N * 10 ], *tris;
void edge( Edge a, Edge b ){
  if(a.tri) a.tri->edge[a.side] = b;
  if(b.tri) b.tri->edge[b.side] = a;
struct Trig { // Triangulation
  Trig(){
    the_root = // Tri should at least contain all
         points
       new(tris++)Tri(Pt(-inf,-inf),Pt(+inf+inf,-inf),Pt
            (-inf,+inf+inf));
  TriRef find(Pt p)const{ return find(the_root,p); }
  void add_point(const Pt& p){ add_point(find(the_root,
  p),p); }
TriRef the_root;
  static TriRef find(TriRef root, const Pt& p) {
    while( true ){
       if( !root->has_chd() )
         return root:
       for( int i = 0; i < 3 \&\& root->chd[i]; ++i)
         if (root->chd[i]->contains(p)) {
           root = root->chd[i];
           break;
    assert( false ); // "point not found"
  void add_point(TriRef root, Pt const& p) {
    TriRef tab, tbc, tca;
      * split it into three triangles */
    tab=new(tris++) Tri(root->p[0],root->p[1],p);
tbc=new(tris++) Tri(root->p[1],root->p[2],p);
    tca=new(tris++) Tri(root->p[2],root->p[0],p);
```

```
edge(Edge(tab,0), Edge(tbc,1));
edge(Edge(tbc,0), Edge(tca,1));
     edge(Edge(tca,0), Edge(tab,1))
     edge(Edge(tab,2), root->edge[2]);
     edge(Edge(tbc,2), root->edge[0]);
     edge(Edge(tca,2), root->edge[1]);
     root->chd[0] = tab;
     root->chd[1] = tbc;
     root->chd[2] = tca;
     flip(tab,2);
     flip(tbc,2);
     flip(tca,2);
  void flip(TriRef tri, SdRef pi) {
    TriRef trj = tri->edge[pi].tri;
     int pj = tri->edge[pi].side;
     if (!trj) return;
if (!in_cc(tri->p[0],tri->p[1],tri->p[2],trj->p[p]
          ])) return;
     /* flip edge between tri,trj */
    TriRef trk = new(tris++) Tri(tri->p[(pi+1)%3], trj
          ->p[pj], tri->p[pi]);
     TriRef trl = new(tris++) Tri(trj->p[(pj+1)%3], tri
          ->p[pi], trj->p[pj]);
     edge(Edge(trk,0), Edge(trl,0));
    edge(Edge(trk,1), tri->edge[(pi+2)%3]);
edge(Edge(trk,2), trj->edge[(pj+1)%3]);
edge(Edge(trl,1), trj->edge[(pj+2)%3]);
    edge(Edge(trl,2), tri->edge[(pi+1)%3]);
tri->chd[0]=trk; tri->chd[1]=trl; tri->chd[2]=0;
trj->chd[0]=trk; trj->chd[1]=trl; trj->chd[2]=0;
     flip(trk,1); flip(trk,2);
     flip(trl,1); flip(trl,2);
  }
vector<TriRef> triang;
set<TriRef> vst;
void go( TriRef now ){
  if( vst.find( now ) != vst.end() )
     return
  vst.insert( now );
  if( !now->has_chd() ){
     triang.push_back( now );
     return:
  for( int i = 0 ; i < now->num_chd() ; i ++ )
     go( now->chd['i ] );
void build( int n , Pt* ps ){
  tris = pool; triang.clear(); vst.clear();
  random\_shuffle(ps, ps + n);
  Trig tri;
  for(int i = 0; i < n; ++ i)
     tri.add_point(ps[i]);
  go( tri.the_root );
4.17 Min Enclosing Circle
```

```
struct Mec{
  // return pair of center and r
  static const int N = 101010;
  int n;
 Pt p[ N ], cen;
  double r2
  void init( int _n , Pt _p[] ){
    n = _n;
    memcpy( p , _p , sizeof(Pt) * n );
  double sqr(double a){ return a*a; }
 Pt center(Pt p0, Pt p1, Pt p2) {
    Pt a = p1-p0;
    Pt b = p2-p0;
    double c1=norm2( a ) * 0.5;
double c2=norm2( b ) * 0.5;
    double d = a \wedge b;
    double x = p0.X + (c1 * b.Y - c2 * a.Y) / d;
    double y = p0.Y + (a.X * c2 - b.X * c1) / d;
    return Pt(x,y);
 pair<Pt,double> solve(){
    random_shuffle(p,p+n);
```

```
r2=0;
       for (int i=0; i<n; i++){</pre>
          if (norm2(cen-p[i]) <= r2) continue;</pre>
          cen = p[i];
          r2 = 0;
          for (int j=0; j<i; j++){
  if (norm2(cen-p[j]) <= r2) continue;
  cen=Pt((p[i].X+p[j].X)/2,(p[i].Y+p[j].Y)/2);</pre>
             r2 = norm2(cen-p[j]);
             for (int k=0; k<j; k++){
  if (norm2(cen-p[k]) <= r2) continue;
  cen = center(p[i],p[k]);</pre>
                r2 = norm2(cen-p[k]);
         }
      return {cen,sqrt(r2)};
} mec;
```

```
4.18 Min Enclosing Ball
// Pt : { x , y , z }
#define N 202020
int n, nouter; Pt pt[ N ], outer[4], res;
double radius,tmp;
double det(double m[3][3]){
     return m[0][0]*m[1][1]*m[2][2]
+ m[0][1]*m[1][2]*m[2][0]
           + m[0][2]*m[2][1]*m[1][0]
- m[0][2]*m[1][1]*m[2][0]
            - m[0][1]*m[1][0]*m[2][2]
            - m[0][0]*m[1][2]*m[2][1];
void ball() {
  Pt q[3]; double m[3][3], sol[3], L[3], d;
  int i, j; res.x = res.y = res.z = radius = 0;
switch ( nouter ) {
     case 1: res=outer[0]; break;
     case 2: res=(outer[0]+outer[1])/2; radius=norm2(res
            outer[0]); break;
     case 3:
       for(i=0; i<2; ++i) q[i]=outer[i+1]-outer[0];
for(i=0; i<2; ++i) for(j=0; j<2; ++j) m[i][j]=(q[i] * q[j])*2;
for(i=0; i<2; ++i) sol[i]=(q[i] * q[i]);
if(fabs(d=m[0][0]*m[1][1]-m[0][1]*m[1][0])<eps)</pre>
       L[0]=(sol[0]*m[1][1]-sol[1]*m[0][1])/d;
L[1]=(sol[1]*m[0][0]-sol[0]*m[1][0])/d;
res=outer[0]+q[0]*L[0]+q[1]*L[1];
        radius=norm2(res, outer[0]);
        break;
     case 4:
        for(i=0; i<3; ++i) q[i]=outer[i+1]-outer[0], sol[
    i]=(q[i] * q[i]);</pre>
        for(i=0;i<3;++i) for(j=0;j<3;++j) m[i][j]=(q[i] *
              q[j])*2;
        d=det(m);
        if(fabs(d)<eps) return;</pre>
        for(j=0; j<3; ++j) {
  for(i=0; i<3; ++i) m[i][j]=sol[i];
  L[j]=det(m) / d;</pre>
           for(i=0; i<3; ++i) m[i][j]=(q[i] * q[j])*2;
        } res=outer[0];
        for(i=0; i<3; ++i ) res = res + q[i] * L[i];</pre>
        radius=norm2(res, outer[0]);
void minball(int n){ ball();
  if(nouter < 4) for(int i = 0; i < n; i ++)
     if(norm2(res, pt[i]) - radius > eps){
        outer[nouter ++] = pt[i]; minball(i); --nouter;
        if(i>0){ Pt Tt = pt[i]
          memmove(&pt[1], &pt[0], sizeof(Pt)*i); pt[0]=Tt
}}}
double solve(){
  // n points in pt
random_shuffle(pt, pt+n); radius=-1;
  for(int i=0;i<n;i++) if(norm2(res,pt[i])-radius>eps)
     nouter=1, outer[0]=pt[i], minball(i);
  return sqrt(radius);
```

## 4.19 Minkowski sum

|}

```
vector<Pt> minkowski(vector<Pt> p, vector<Pt> q){
  int n = p.size() , m = q.size();
  Pt c = Pt(0, 0);
  for( int i = 0; i < m; i ++) c = c + q[i];
  c = c / m;
  for( int i = 0; i < m; i ++) q[i] = q[i] - c;
  int cur = -1;
  for( int i = 0; i < m; i ++)</pre>
    if( (q[i] ^ (p[0] - p[n-1])) >
                                     -eps)
      if( cur == -1 || (q[i] \wedge (p[0] - p[n-1])) >
                        (q[cur] ^ (p[0] - p[n-1])) )
        cur = i:
  vector<Pt> h;
  p.push_back(p[0]);
  for( int i = 0; i < n; i ++)
  while( true ){</pre>
      h.push_back(p[i] + q[cur]);
      int nxt = (cur + 1 == m ? 0 : cur + 1);
if((q[cur] ^ (p[i+1] - p[i])) < -eps) cur = nxt;
      else break;
    }
  for(auto &&i : h) i = i + c;
  return convex_hull(h);
```

## 4.20 Min dist on Cuboid

```
typedef LL T;
Tr;
if(i>=0 && i< 2) turn(i+1, j, x0+L+z, y, x0+L-x, x0+L, y0, H, W, L);
  if(j>=0 && j< 2) turn(i, j+1, x, y0+W+z, y0+W-y, x0, y0+W, L, H, W);
 if(i<=0 && i>-2) turn(i-1, j, x0-z, y, x-x0, x0-H, y0, H, W, L);
if(j<=0 && j>-2) turn(i, j-1, x, y0-z, y-y0,
                           x0, y0-H, L, H, W);
Ť solve(T L, T W, T H,
         T x1, T y1, T z1, T x2, T y2, T z2){
  if( z1!=0 && z1!=H ){
    if( y1==0 || y1==W )
      swap(y1,z1), swap(y2,z2), swap(W,H);
    else swap(x1,z1), swap(x2,z2), swap(L,H);
  if (z1==H) z1=0, z2=H-z2;
  r=INF; turn(0,0,x2-x1,y2-y1,z2,-x1,-y1,L,W,H);
  return r;
```

## 4.21 Heart of Triangle

```
Pt inCenter( Pt &A,
                      Pt &B, Pt &C) { // 内心
  double a = norm(B-C), b = norm(C-A), c = norm(A-B);
  return (A * a + B * b + C * c) / (a + b + c);
Pt circumCenter( Pt &a, Pt &b, Pt &c) { // 外心
  Pt bb = b - a, cc = c - a;
  double db=norm2(bb), dc=norm2(cc), d=2*(bb ^ cc);
  return a-Pt(bb.Y*dc-cc.Y*db, cc.X*db-bb.X*dc) / d;
Pt othroCenter( Pt &a, Pt &b, Pt &c) { // {\scriptstyle{\#}} {\scriptstyle{\psi}}
  Pt ba = b - a, ca = c - a, bc = b - c;
double Y = ba.Y * ca.Y * bc.Y,
    A = ca.X * ba.Y - ba.X * ca.Y
    x0= (Y+ca.X*ba.Y*b.X-ba.X*ca.Y*c.X) / A,
    y0 = -ba.X * (x0 - c.X) / ba.Y + ca.Y;
  return Pt(x0, y0);
```

# Graph

#### 5.1 DominatorTree

```
const int MAXN = 100010;
struct DominatorTree{
#define REP(i,s,e) for(int i=(s);i<=(e);i++)</pre>
#define REPD(i,s,e) for(int i=(s);i>=(e);i--)
  int n , m , s;
vector< int > g[ MAXN ] , pred[ MAXN ];
vector< int > cov[ MAXN ];
int dfn[ MAXN ] , nfd[ MAXN ] , ts;
  int par[ MAXN ];
  int sdom[ MAXN ] , idom[ MAXN ];
int mom[ MAXN ] , mn[ MAXN ];
inline bool cmp( int u , int v )
   { return dfn[ u ] < dfn[ v ]; }
   int eval( int u ){
     if( mom[ u ] == u ) return u;
     int res = eval( mom[ u ] );
     if(cmp( sdom[ mn[ mom[ u j j ] ] , sdom[ mn[ u ] ] ))
   mn[ u ] = mn[ mom[ u ] ];
     return mom[ u ] = res;
  void init( int _n , int _m , int _s ){
     ts = 0; n = _n; m = _m; s = _s;
     REP( i, 1, n ) g[ i ].clear(), pred[ i ].clear();
  void addEdge( int u , int v ){
     g[ u ].push_back( v );
pred[ v ].push_back( u );
  void dfs( int u ){
     ts++;
     dfn['u ] = ts;
     nfd[ts] = u;
     for( int v : g[ u ] ) if( dfn[ v ] == 0 ){
  par[ v ] = u;
        dfs(v);
  void build(){
     REP( i , 1 , n ){
    dfn[ i ] = nfd[ i ] = 0;
    cov[ i ].clear();
       mom[i] = mn[i] = sdom[i] = i;
     dfs( s );
     REPD( i , n , 2 ){
  int u = nfd[ i ];
        if( u == 0 ) continue :
        for( int v : pred[ u ] ) if( dfn[ v ] ){
          eval( v );
          if( cmp( sdom[ mn[ v ] ] , sdom[ u ] ) )
             sdom[u] = sdom[mn[v]];
       cov[ sdom[ u ] ].push_back( u );
       mom[ u ] = par[ u ];
       for( int w : cov[ par[ u ] ] ){
          eval( w );
          if( cmp( sdom[ mn[ w ] ] , par[ u ] ) )
          idom[w] = mn[w];
else idom[w] = par[u];
       cov[ par[ u ] ].clear();
     REP( i , 2 , n ){
  int u = nfd[ i ];
        if( u == 0 ) continue ;
        if( idom[ u ] != sdom[ u ] )
          idom[ u ] = idom[ idom[ u ] ];
  }
} domT;
5.2 MaxClique
```

```
#define N 111
struct MaxClique{ // 0-base
  typedef bitset< N > Int;
  Int linkto[N], v[N];
  int n;
  void init( int _n ){
    for( int i = 0 ; i < n ; i ++ ){</pre>
      linkto[ i ].reset();
```

```
v[ i ].reset();
    }
  void addEdge( int a , int b ){
    v[a][b] = v[b][a] = 1;
  int popcount(const Int& val)
  { return val.count(); }
  int lowbit(const Int& val)
  { return val._Find_first(); } int ans , stk[ N ]; int id[ N ] , di[ N ] , deg[ N ];
  Int cans;
  void maxclique(int elem_num, Int candi){
    if(elem_num > ans){
       ans = elem_num;
       cans.reset();
for( int i = 0 ; i < elem_num ; i ++ )</pre>
         cans[ id[ stk[ i ] ] = 1;
     int potential = elem_num + popcount(candi);
     if(potential <= ans) return;</pre>
     int pivot = lowbit(candi);
     Int smaller_candi = candi & (~linkto[pivot]);
    while(smaller_candi.count() && potential > ans){
       int next = lowbit(smaller_candi);
       candi[next] = !candi[next];
       smaller_candi[ next ] = !smaller_candi[ next ];
       potential -
       if(next == pivot || (smaller_candi & linkto[next
            ]).count() ){
         stk[elem_num] = next;
         maxclique(elem_num + 1, candi & linkto[next]);
    }
  int solve(){
     for( int i = 0 ; i < n ; i ++ ){</pre>
       id[ i ] = i;
deg[ i ] = v[ i ].count();
     sort( id , id + n , [&](int id1, int id2){
    return deg[id1] > deg[id2]; } );
     for( int i = 0 ; i < n ; i ++ )
       di[ id[ i ] ] = i;
    for( int i = 0 ; i < n ; i ++ )
  for( int j = 0 ; j < n ; j ++ )
    if( v[ i ][ j ] )
      linkto[ di[ i ] ][ di[ j ] ] = 1;</pre>
     Int cand; cand.reset();
     for( int i = 0 ; i < n ; i ++ )</pre>
       cand[i] = 1;
     ans = 1;
    cans.reset(); cans[0] = 1;
    maxclique(0, cand);
     return ans;
} solver;
5.3 MaxCliqueDyn
```

```
#define N 150
struct MaxClique{ // Maximum Clique
 bitset<N> a[N],cs[N];
  int ans,sol[N],q,cur[N],d[N],n;
 void init(int _n){
   n=_n; for(int i=0;i<n;i++) a[i].reset();
  void addEdge(int u,int v){a[u][v]=a[v][u]=1;}
 void csort(vector<int> &r,vector<int> &c){
   int mx=1,km=max(ans-q+1,1),t=0,m=r.size();
    cs[1].reset(); cs[2].reset();
    for(int i=0;i<m;i++){</pre>
      int p=r[i],k=1;
      while((cs[k]&a[p]).count()) k++;
      if(k>mx){ mx++; cs[mx+1].reset();}
      cs[k][p]=1;
      if(k<km) r[t++]=p;
   }
   c.resize(m);
    if(t) c[t-1]=0;
    for(int k=km;k<=mx;k++){</pre>
```

```
for(int p=cs[k]._Find_first();p<N;p=cs[k].</pre>
            Find_next(p)){
         r[t]=p; c[t]=k; t++;
    }
  }
  void dfs(vector<int> &r,vector<int> &c,int l,bitset<N</pre>
       > mask){
    while(!r.empty()){
      int p=r.back(); r.pop_back(); mask[p]=0;
       if(q+c.back()<=ans) return;</pre>
       cur[q++]=p;
       vector<int> nr,nc; bitset<N> nmask=mask&a[p];
       for(int i:r) if(a[p][i]) nr.push_back(i);
       if(!nr.empty()){
         if(1<4){
           for(int i:nr) d[i]=(a[i]&nmask).count();
sort(nr.begin(),nr.end(),[&](int x,int y){
                return d[x]>d[y];});
         csort(nr,nc); dfs(nr,nc,l+1,nmask);
      else if(q>ans){
        ans=q; copy(cur,cur+q,sol);
       c.pop_back(); q--;
  }
  int solve(bitset<N> mask){ // vertex mask
    vector<int> r,c;
    for(int i=0;i<n;i++) if(mask[i]) r.push_back(i);</pre>
    for(int i=0;i<n;i++) d[i]=(a[i]&mask).count();</pre>
    sort(r.begin(),r.end(),[&](int i,int j){return d[i
         ]>d[j];})
    csort(r,c); dfs(r,c,1,mask);
    return ans; // sol[0 ~ ans-1]
}graph;
```

## 5.4 Strongly Connected Component

```
void dfs(int i){
  V[i]=low[i]=++ts,stk[top++]=i,instk[i]=1;
  for(auto x:E[i]){
    if(!V[x])dfs(x),low[i]=min(low[i],low[x]);
    else if(instk[x])low[i]=min(low[i],V[x]);
  if(V[i]==low[i]){
    int j;
do{j = stk[--top], instk[j] = 0, scc[j] = i;
    }while(j != i);
  }
}
```

## 5.5 Dynamic MST

```
/* Dynamic MST 0( Q lg^2 Q )
 n nodes, m edges, Q query
 (u[i], v[i], w[i])->edge
 (qid[i], qw[i])->chg weight of edge No.qid[i] to qw[i]
 delete an edge: (i, \infty)
add an edge: change from \infty to specific value */
const int SZ=M+3*MX0:
int a[N],*tz;
int find(int x){
     return x==a[x]?x:a[x]=find(a[x]);
bool cmp(int aa,int bb){ return tz[aa]<tz[bb]; }
int kx[N],ky[N],kt, vd[N],id[M], app[M], cur;
long long answer[MXQ]; // answer after ith query</pre>
bool extra[M];
void solve(int *qx,int *qy,int Q,int n,int *x,int *y,
      int *z, int m1, long long ans){
  if(Q==1){
     for(int i=1;i<=n;i++) a[i]=0;</pre>
     z[qx[0]]=qy[0]; tz = z;
for(int i=0;i<m1;i++) id[i]=i;
     sort(id,id+m1,cmp); int ri,rj;
for(int i=0;i<m1;i++){</pre>
        ri=find(x[id[i]]); rj=find(y[id[i]]);
        if(ri!=rj){    ans+=z[id[i]];    a[ri]=rj;    }
```

```
answer[cur++]=ans;
     return;
  int ri,rj;
  //contract
  kt=0;
  for(int i=1;i<=n;i++) a[i]=0;</pre>
  for(int i=0;i<Q;i++){</pre>
     ri=find(x[qx[i]]); rj=find(y[qx[i]]); if(ri!=rj) a[
  int tm=0;
  for(int i=0;i<m1;i++) extra[i]=true;</pre>
  for(int i=0;i<0;i++) extra[ qx[i] ]=false;
for(int i=0;i<m1;i++) if(extra[i]) id[tm++]=i;</pre>
  tz=z; sort(id,id+tm,cmp);
  for(int i=0;i<tm;i++){
    ri=find(x[id[i]]);    rj=find(y[id[i]]);</pre>
     if(ri!=rj){
       a[ri]=rj; ans += z[id[i]];
kx[kt]=x[id[i]]; ky[kt]=y[id[i]]; kt++;
    }
  for(int i=1;i<=n;i++) a[i]=0;
  for(int i=0;i<kt;i++) a[ find(kx[i]) ]=find(ky[i]);</pre>
  int n2=0;
  for(int i=1;i<=n;i++) if(a[i]==0)</pre>
  vd[i]=++n2;
  for(int i=1;i<=n;i++) if(a[i])</pre>
  vd[i]=vd[find(i)];
int m2=0, *Nx=x+m1, *Ny=y+m1, *Nz=z+m1;
  for(int i=0;i<m1;i++) app[i]=-1;</pre>
  for(int i=0;i<0;i++) if(app[qx[i]]==-1){
  Nx[m2]=vd[ x[ qx[i] ] ];  Ny[m2]=vd[ y[ qx[i] ] ];
        Nz[m2]=z[ qx[i] ];</pre>
     app[qx[i]]=m2; m2++;
  for(int i=0;i<Q;i++){ z[ qx[i] ]=qy[i]; qx[i]=app[qx[</pre>
  i]]; }
for(int i=1;i<=n2;i++) a[i]=0;
  for(int i=0;i<tm;i++){</pre>
     ri=find(vd[ x[id[i]] ]); rj=find(vd[ y[id[i]] ]);
     if(ri!=rj){
       a[ri]=rj; Nx[m2]=vd[ x[id[i]] ];
       Ny[m2]=vd[ y[id[i]] ]; Nz[m2]=z[id[i]]; m2++;
    }
  int mid=Q/2;
  solve(qx,qy,mid,n2,Nx,Ny,Nz,m2,ans);
  solve(qx+mid,qy+mid,Q-mid,n2,Nx,Ny,Nz,m2,ans);
int u[SZ],v[SZ],w[SZ],qid[MXQ],qw[MXQ],n,m,Q;
void work(){if(Q) cur=0,solve(qid,qw,Q,n,u,v,w,m,0);}
```

## 5.6 Maximum General graph Matching

```
// should shuffle vertices and edges
const int N = 100005, E = (2e5) * 2 + 40;
struct Graph{
  int to[E],bro[E],head[N],e;
  int lnk[N],vis[N],stp,n;
void init( int _n ){
   stp = 0; e = 1; n = _n;
     for( int i = 1 ; i <= n ; i ++ )</pre>
       lnk[i] = vis[i] = 0;
  void add_edge(int u,int v){
    to[e]=v,bro[e]=head[u],head[u]=e++;
to[e]=u,bro[e]=head[v],head[v]=e++;
  bool dfs(int x){
    vis[x]=stp;
    for(int i=head[x];i;i=bro[i]){
       int v=to[i];
       if(!lnk[v]){
         lnk[x]=v, lnk[v]=x;
         return true;
    } for(int i=head[x];i;i=bro[i]){
       int v=to[i];
       if(vis[lnk[v]]<stp){</pre>
         int w=lnk[v];
```

```
lnk[x]=v,lnk[v]=x,lnk[w]=0;
    if(dfs(w)) return true;
    lnk[w]=v,lnk[v]=w,lnk[x]=0;
}
return false;
}
int solve(){
  int ans = 0;
  for(int i=1;i<=n;i++) if(!lnk[i])
    stp++, ans += dfs(i);
  return ans;
}
} graph;</pre>
```

## 5.7 Minimum General Weighted Matching

```
struct Graph {
  // Minimum General Weighted Matching (Perfect Match)
  static const int MXN = 105;
  int n, edge[MXN][MXN];
  int match[MXN],dis[MXN],onstk[MXN];
  vector<int> stk;
  void init(int _n) {
    n = _n;
for( int i = 0 ; i < n ; i ++ )</pre>
       for( int j = 0 ; j < n ; j ++ )
  edge[ i ][ j ] = 0;</pre>
  void add_edge(int u, int v, int w)
  \{ edge[u][v] = edge[v][u] = w; \}
  bool SPFA(int u){
    if (onstk[u]) return true;
    stk.PB(u);
    onstk[u] = 1;
    for (int v=0; v<n; v++){</pre>
       if (u != v && match[u] != v && !onstk[v]){
         int m = match[v];
         if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
            dis[m] = dis[u] - edge[v][m] + edge[u][v];
            onstk[v] = 1;
            stk.PB(v)
            if (SPFA(m)) return true;
            stk.pop_back();
            onstk[v] = 0;
       }
    onstk[u] = 0
    stk.pop_back();
    return false;
  int solve() {
    // find a match
    for (int i=0; i<n; i+=2){</pre>
       match[i] = i+1;
       match[i+1] = i;
    while (true){
       int found = 0;
       for( int i = 0 ; i < n ; i ++ )
  onstk[ i ] = dis[ i ] = 0;
for (int i=0; i<n; i++){</pre>
         stk.clear()
         if (!onstk[i] && SPFA(i)){
            found = 1
            while (SZ(stk)>=2){
              int u = stk.back(); stk.pop_back();
int v = stk.back(); stk.pop_back();
              match[u] = v;
              match[v] = u;
         }
       if (!found) break;
    int ret = 0;
    for (int i=0; i<n; i++)
      ret += edge[i][match[i]];
    ret /= 2;
    return ret;
```

|}graph;

## 5.8 Maximum General Weighted Matching

```
struct WeightGraph {
  static const int INF = INT_MAX;
  static const int N = 514;
  struct edge{
    int u,v,w; edge(){}
    edge(int ui,int vi,int wi)
      :u(ui),v(vi),w(wi){}
  int n,n_x
 edge g[N*2][N*2];
  int lab[N*2]
  int match[N*2],slack[N*2],st[N*2],pa[N*2];
  int flo_from[N*2][N+1],S[N*2],vis[N*2];
 vector<int> flo[N*2];
 queue<int> q;
  int e_delta(const edge &e){
   return lab[e.u]+lab[e.v]-g[e.u][e.v].w*2;
 void update_slack(int u,int x){
    if(!slack[x]||e_delta(g[u][x])<e_delta(g[slack[x]][</pre>
        x]))slack[x]=u;
  void set_slack(int x){
    slack[x]=0;
    for(int u=1;u<=n;++u)</pre>
      if(g[u][x].w>0&&st[u]!=x&&S[st[u]]==0)
        update_slack(u,x);
 void q_push(int x){
    if(x<=n)q.push(x);</pre>
    else for(size_t i=0;i<flo[x].size();i++)</pre>
      q_push(flo[x][i]);
  void set_st(int x,int b){
    st[x]=b;
    if(x>n)for(size_t i=0;i<flo[x].size();++i)</pre>
      set_st(flo[x][i],b);
  int get_pr(int b,int xr){
    int pr=find(flo[b].begin(),flo[b].end(),xr)-flo[b].
        begin();
    if(pr%2==1){
      reverse(flo[b].begin()+1,flo[b].end());
      return (int)flo[b].size()-pr;
   }else return pr;
  void set_match(int u,int v){
   match[u]=g[u][v].v;
    if(u<=n) return;</pre>
    edge e=g[u][v];
    int xr=flo_from[u][e.u],pr=get_pr(u,xr);
    for(int i=0;i<pr;++i)set_match(flo[u][i],flo[u][i</pre>
        ^17);
    set_match(xr,v);
    rotate(flo[u].begin(),flo[u].begin()+pr,flo[u].end
  void augment(int u,int v){
    for(;;){
      int xnv=st[match[u]];
      set_match(u,v);
      if(!xnv)return;
      set_match(xnv,st[pa[xnv]]);
      u=st[pa[xnv]],v=xnv;
  int get_lca(int u,int v){
    static int t=0;
    for(++t;ullv;swap(u,v)){
      if(u==0)continue;
      if(vis[u]==t)return u;
      vis[u]=t;
      u=st[match[u]];
      if(u)u=st[pa[u]];
   }
    return 0;
  void add_blossom(int u,int lca,int v){
```

```
int b=n+1:
  while(b<=n_x&&st[b])++b;</pre>
  if(b>n_x)++n_x
  lab[b]=0,S[b]=0;
  match[b]=match[lca];
  flo[b].clear();
  flo[b].push_back(lca);
  for(int x=u,y;x!=lca;x=st[pa[y]])
    flo[b].push_back(x),flo[b].push_back(y=st[match[x
         ]]),q_push(y)
  reverse(flo[b].begin()+1,flo[b].end());
  for(int x=v,y;x!=lca;x=st[pa[y]])
    flo[b].push\_back(x),flo[b].push\_back(y=st[match[x
         ]]),q_push(y);
  set_st(b,b);
  for(int x=1;x<=n_x;++x)g[b][x].w=g[x][b].w=0;
  for(int x=1;x<=n;++x)flo_from[b][x]=0;
for(size_t i=0;i<flo[b].size();++i){</pre>
    int xs=flo[b][i];
    for(int x=1;x<=n_x;++x)</pre>
       if(g[b][x].w==0|le_delta(g[xs][x])<e_delta(g[b][x])
            ][x]))
         g[b][x]=g[xs][x],g[x][b]=g[x][xs];
    for(int x=1;x<=n;++x)</pre>
       if(flo_from[xs][x])flo_from[b][x]=xs;
  set_slack(b);
}
void expand_blossom(int b){
  for(size_t i=0;i<flo[b].size();++i)
  set_st(flo[b][i],flo[b][i]);</pre>
  int xr=flo_from[b][g[b][pa[b]].u],pr=get_pr(b,xr);
  for(int i=0;i<pr;i+=2){
  int xs=flo[b][i],xns=flo[b][i+1];</pre>
    pa[xs]=g[xns][xs].u;
    S[xs]=1,S[xns]=0;
    slack[xs]=0,set_slack(xns);
    q_push(xns);
  S[xr]=1,pa[xr]=pa[b];
  for(size_t i=pr+1;i<flo[b].size();++i){</pre>
    int xs=flo[b][i]
    S[xs]=-1, set\_slack(xs);
  st[b]=0;
bool on_found_edge(const edge &e){
  int u=st[e.u],v=st[e.v];
  if(S[v]==-1)
    pa[v]=e.u,S[v]=1;
    int nu=st[match[v]];
    slack[v]=slack[nu]=0;
  S[nu]=0,q_push(nu);
}else if(S[v]==0){
    int lca=get_lca(u,v);
    if(!lca)return augment(u,v),augment(v,u),true;
    else add_blossom(u,lca,v);
  return false;
bool matching(){
  memset(S+1,-1,sizeof(int)*n_x);
  memset(slack+1,0,sizeof(int)*n_x);
  q=queue<int>();
  for(int x=1;x<=n_x;++x)</pre>
    if(st[x]==x\&\&!match[x])pa[x]=0,S[x]=0,q_push(x);
  if(q.empty())return false;
  for(;;){
  while(q.size()){
       int u=q.front();q.pop();
       if(S[st[u]]==1)continue;
       for(int v=1;v<=n;++v)</pre>
         if(g[u][v].w>0&&st[u]!=st[v]){
           if(e_delta(g[u][v])==0){
  if(on_found_edge(g[u][v]))return true;
           }else update_slack(u,st[v]);
         }
    int d=INF;
    for(int b=n+1;b<=n_x;++b)</pre>
       if(st[b]==b&&S[b]==1)d=min(d,lab[b]/2);
    for(int x=1;x<=n_x;++x)</pre>
```

```
if(st[x]==x\&slack[x]){
           if(S[x]==-1)d=min(d,e_delta(g[slack[x]][x]))
           else if(S[x]==0)d=min(d,e_delta(g[slack[x]][x
               1)/2);
      for(int u=1;u<=n;++u){</pre>
        if(S[st[u]]==0){
           if(lab[u]<=d)return 0;</pre>
           lab[u]-=d;
        }else if(S[st[u]]==1)lab[u]+=d;
      for(int b=n+1;b<=n_x;++b)</pre>
        if(st[b]==b){
          if(S[st[b]]==0)lab[b]+=d*2;
           else if(S[st[b]]==1)lab[b]-=d*2;
      q=queue<int>();
      for(int x=1;x<=n_x;++x)</pre>
        if(st[x]==x&&slack[x]&&st[slack[x]]!=x&&e_delta
             (g[slack[x]][x])==0)
           if(on_found_edge(g[slack[x]][x]))return true;
      for(int b=n+1;b<=n_x;++b)</pre>
        if(st[b]==b\&\&S[b]==1\&\&lab[b]==0)expand_blossom(
             b);
    return false;
  pair<long long,int> solve(){
    memset(match+1,0,sizeof(int)*n);
    int n_matches=0;
    long long tot_weight=0;
    for(int u=0;u<=n;++u)st[u]=u,flo[u].clear();</pre>
    int w_max=0;
    for(int u=1;u<=n;++u)</pre>
      for(int v=1;v<=n;++v){</pre>
        flo_from[u][v]=(u==v?u:0);
        w_max=max(w_max,g[u][v].w);
    for(int u=1;u<=n;++u)lab[u]=w_max;</pre>
    while(matching())++n_matches;
    for(int u=1;u<=n;++u)</pre>
      if(match[u]&&match[u]<u)</pre>
        tot_weight+=g[u][match[u]].w;
    return make_pair(tot_weight,n_matches);
  void add_edge( int ui , int vi , int wi ){
    g[ui][vi].w = g[vi][ui].w = wi;
  void init( int _n ){
    n = _n;
    for(int u=1;u<=n;++u)</pre>
      for(int v=1;v<=n;++v)</pre>
        g[u][v]=edge(u,v,0);
} graph;
```

# 5.9 Minimum Steiner Tree

```
if( i != j && dst[ i ][ j ] != INF )
  dst[ i ][ j ] += w[ i ];
for( int k = 0 ; k < n ; k ++ )</pre>
         for( int i = 0 ; i < n ; i ++ )</pre>
            for( int i = 0 ; i < n ; i ++
for( int j = 0 ; j < n ; j ++
  if( dst[ i ][ j ] != INF )
    dst[ i ][ j ] += w[ j ];</pre>
   int solve( const vector<int>& ter ){
      int t = (int)ter.size();
for( int i = 0 ; i < (1 << t ) ; i ++ )</pre>
      for(int j = 0; j < n; j ++)

dp[i][j] = INF;

for(int i = 0; i < n; i ++)
         d\hat{p}[0][i] = 0;
      for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){
  if( msk == ( msk & (-msk) ) ){</pre>
            int who = __lg( msk );
for( int i = 0 ; i < n ; i ++ )
  dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];</pre>
            continue;
         for( int i = 0 ; i < n ; i ++ )
            for( int submsk = ( msk - 1 ) & msk ; submsk ;
     submsk = ( submsk - 1 ) & msk )
                  [i]
for(int i = 0; i < n; i ++){
                                                          );
            tdst[ i ] = INF;
               or( int j = 0 ; j < n ; j ++ )
tdst[ i ] = min( tdst[ i ],
            for( int j = 0;
                                dp[ msk ][ j ] + dst[ j ][ i ] - w
                                       [j]);
         for( int i = 0 ; i < n ; i ++ )
  dp[ msk ][ i ] = tdst[ i ];</pre>
      int ans = INF;
      for( int i = 0 ; i < n ; i ++ )
  ans = min( ans , dp[ ( 1 << t ) - 1 ][ i ] );</pre>
      return ans:
} solver;
```

#### 5.10 BCC based on vertex

```
struct BccVertex {
  int n,nScc,step,dfn[MXN],low[MXN];
  vector<int> E[MXN],sccv[MXN];
  int top,stk[MXN];
  void init(int _n) {
    n = _n; nScc = step = 0;
for (int i=0; i<n; i++) E[i].clear();</pre>
  void addEdge(int u, int v)
  { E[u].PB(v); E[v].PB(u); } void DFS(int u, int f) {
    dfn[u] = low[u] = step++;
    stk[top++] = u
    for (auto v:E[u]) {
       if (v == f) continue;
       if (dfn[v] == -1) {
         DFS(v,u);
         low[u] = min(low[u], low[v]);
         if (low[v] >= dfn[u]) {
           sccv[nScc].clear();
           do {
             z = stk[--top];
             sccv[nScc].PB(z);
           } while (z != v)
           sccv[nScc++].PB(u);
       }else
         low[u] = min(low[u],dfn[v]);
```

```
#define INF 1000000000000000LL
  vector<vector<int>> solve() {
                                                                   #define N 5010
    vector<vector<int>> res;
                                                                   #define M 200010
     for (int i=0; i<n; i++)</pre>
                                                                   struct edge{
       dfn[i] = low[i] = -1;
                                                                     int to; LL w;
     for (int i=0; i<n; i++)
                                                                     edge(int a=0, LL b=0): to(a), w(b){}
       if (dfn[i] == -1) {
         top = 0;
                                                                   struct node{
         DFS(i,i);
                                                                     LL d; int u, next;
                                                                     node(LL a=0, int b=0, int c=0): d(a), u(b), next(c){}
    REP(i,nScc) res.PB(sccv[i]);
                                                                   struct DirectedGraphMinCycle{
    return res;
                                                                     vector<edge> g[N], grev[N];
}graph;
                                                                     LL dp[N][N], p[N], d[N], mu;
                                                                     bool inq[N];
5.11 Min Mean Cycle
                                                                     int n, bn, bsz, hd[N];
                                                                     void b_insert(LL d, int u){
                                                                        int i = d/mu;
/* minimum mean cycle O(VE) */
struct MMC{
                                                                        if(i >= bn) return;
#define E 101010
                                                                        b[++bsz] = node(d, u, hd[i]);
#define V 1021
                                                                        hd[i] = bsz;
#define inf 1e9
                                                                     void init( int _n ){
#define eps 1e-6
                                                                       n = _n;
for( int i = 1 ; i <= n ; i ++ )
  g[ i ].clear();
  struct Edge { int v,u; double c; };
  int n, m, prv[V][V], prve[V][V], vst[V];
  Edge e[E];
  vector<int> edgeID, cycle, rho;
  double d[V][V];
                                                                     void addEdge( int ai , int bi , LL ci )
  void init( int _n )
                                                                      { g[ai].push_back(edge(bi,ci)); }
  { n = _n; m = 0; }
// WARNING: TYPE matters
                                                                        solve(){
                                                                        fill(dp[0], dp[0]+n+1, 0);
  void addEdge( int vi , int ui , double ci )
{ e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
                                                                        for(int i=1; i<=n; i++){
                                                                          fill(dp[i]+1, dp[i]+n+1, INF);
for(int j=1; j<=n; j++) if(dp[i-1][j] < INF){
   for(int k=0; k<(int)g[j].size(); k++)</pre>
    dp[i][g[j][k].to] =min(dp[i][g[j][k].to]
                                                                                                         dp[i-1][j]+g[j][k].w);
                                                                          }
                                                                        }
                                                                        mu=INF; LL bunbo=1;
           d[i+1][u] = d[i][v]+e[j].c;
                                                                        for(int i=1; i<=n; i++) if(dp[n][i] < INF){</pre>
           prv[i+1][u] = v;
                                                                          LL a=-INF, b=1;
           prve[i+1][u] = j;
                                                                          for(int j=0; j<=n-1; j++) if(dp[j][i] < INF){
   if(a*(n-j) < b*(dp[n][i]-dp[j][i])){</pre>
         }
                                                                               a = dp[n][i]-dp[j][i];
      }
    }
                                                                               b = n-j;
                                                                            }
  double solve(){
    // returns inf if no cycle, mmc otherwise
                                                                          if(mu*b > bunbo*a)
    double mmc=inf;
                                                                            mu = a, bunbo = b;
    int st = -1;
    bellman_ford();
                                                                        if(mu < 0) return -1; // negative cycle</pre>
                                                                        if(mu == INF) return INF; // no cycle
    for(int i=0; i<n; i++) {
       double avg=-inf;
                                                                        if(mu == 0) return 0;
                                                                        for(int i=1; i<=n; i++)</pre>
       for(int k=0; k<n; k++) {
  if(d[n][i]<inf-eps) avg=max(avg,(d[n][i]-d[k][i])</pre>
                                                                          for(int j=0; j<(int)g[i].size(); j++)</pre>
                                                                          g[i][j].w *= bunbo;
              1)/(n-k);
         else avg=max(avg,inf);
                                                                        memset(p, 0, sizeof(p));
                                                                        queue<int> q;
                                                                        for(int i=1; i<=n; i++){</pre>
       if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
                                                                          q.push(i);
     if(st==-1) return inf;
                                                                          inq[i] = true;
    FZ(vst);edgeID.clear();cycle.clear();rho.clear();
     for (int i=n; !vst[st]; st=prv[i--][st]) {
                                                                        while(!q.empty()){
                                                                          int i=q.front(); q.pop(); inq[i]=false;
for(int j=0; j<(int)g[i].size(); j++){</pre>
       vst[st]++;
       edgeID.PB(prve[i][st]);
       rho.PB(st);
                                                                             if(p[g[i][j].to] > p[i]+g[i][j].w-mu){
                                                                               p[g[i][j].to] = p[i]+g[i][j].w-mu;
if(!inq[g[i][j].to]){
    while (vst[st] != 2) {
       int v = rho.back(); rho.pop_back();
                                                                                 q.push(g[i][j].to);
       cycle.PB(v);
                                                                                 inq[g[i][j].to] = true;
                                                                               }
       vst[v]++;
                                                                            }
                                                                          }
    reverse(ALL(edgeID));
    edgeID.resize(SZ(cycle));
                                                                        for(int i=1; i<=n; i++) grev[i].clear();</pre>
    return mmc:
                                                                        for(int i=1; i<=n; i++)</pre>
                                                                          for(int j=0; j<(int)g[i].size(); j++){
  g[i][j].w += p[i]-p[g[i][j].to];</pre>
} mmc;
```

grev[g[i][j].to].push\_back(edge(i, g[i][j].w));

LL mldc = n\*mu;

## // works in O(N M)

Directed Graph Min Cost Cycle

```
National Taiwan University CRyptoGRapheR
     for(int i=1; i<=n; i++){
  bn=mldc/mu, bsz=0;
  memset(hd, 0, sizeof(hd));</pre>
                                                                           heap* merge(heap* curNd, heap* newNd){
       fill(d+i+1, d+n+1, INF);
b_insert(d[i]=0, i);
        for(int j=0; j<=bn-1; j++) for(int k=hd[j]; k; k=
            b[k].next){
          int u = b[k].u;
          LL du = b[k].d;
          if(du > d[u]) continue;
for(int l=0; l<(int)g[u].size(); l++) if(g[u][l
       ].to > i){
            if(d[g[u][l].to] > du + g[u][l].w){
               d[g[u][l].to] = du + g[u][l].w;
               b_insert(d[g[u][l].to], g[u][l].to);
         }
       for(int j=0; j<(int)grev[i].size(); j++) if(grev[</pre>
             i][j].to > i)
          mldc=min(mldc,d[grev[i][j].to] + grev[i][j].w);
     return mldc / bunbo;
} graph;
5.13 K-th Shortest Path
// time: O(|E| \setminus lg \mid E| + |V| \setminus lg \mid V| + K)
// memory: 0(|E| \lg |E| + |V|)
struct KSP{ // 1-base
  struct nd{
     int u, v, d;
     nd(int ui = 0, int vi = 0, int di = INF)
     \{ u = ui; v = vi; d = di; \}
  };
  struct heap{
    nd* edge; int dep; heap* chd[4];
  static int cmp(heap* a,heap* b)
  { return a->edge->d > b->edge->d; }
  struct node{
     int v; LL d; heap* H; nd* E;
     node(){}
    node(LL _d, int _v, nd* _E) { d =_d; v = _v; E = _E; } node(heap* _H, LL _d)
     \{ H = _H; d = _d; \}
```

friend bool operator<(node a, node b)</pre>

for( int i = 1; i <= n; i ++ ){
 g[ i ].clear(); rg[ i ].clear();
 nxt[ i ] = head[ i ] = NULL;
 dst[ i ] = -1;</pre>

void addEdge( int ui , int vi , int di ){
 nd\* e = new nd(ui, vi, di);
 g[\_ui ].push\_back( e );

void init( int \_n , int \_k , int \_s , int \_t ){
 n = \_n; k = \_k; s = \_s; t = \_t;

{ return a.d > b.d; }

int n, k, s, t, dst[ N ];
nd \*nxt[ N ];

vector<nd\*> g[ N ], rg[ N ];
heap \*nullNd, \*head[ N ];

rg[ vi ].push\_back( e );

priority\_queue<node> Q; Q.push(node(0, t, NULL));

while (!Q.empty()){

dst[ p.v ] = p.d;
nxt[ p.v ] = p.E;
dfsQ.push( p.v );

for(auto e: rg[ p.v ])

while(dfsQ.size()) dfsQ.pop();

node p = Q.top(); Q.pop();
if(dst[p.v] != -1) continue;

Q.push(node(p.d + e->d, e->u, e));

queue<int> dfsQ; void dijkstra(){

```
if(curNd == nullNd) return newNd;
     heap* root = new heap;
     memcpy(root, curNd, sizeof(heap));
     if(newNd->edge->d < curNd->edge->d){
        root->edge = newNd->edge;
root->chd[2] = newNd->chd[2]
        root->chd[3] = newNd->chd[3];
        newNd->edge = curNd->edge;
newNd->chd[2] = curNd->chd[2];
        newNd \rightarrow chd[3] = curNd \rightarrow chd[3];
     if(root->chd[0]->dep < root->chd[1]->dep)
        root->chd[0] = merge(root->chd[0],newNd);
        root->chd[1] = merge(root->chd[1],newNd);
     root->dep = max(root->chd[0]->dep, root->chd[1]->
          dep) + 1;
     return root;
  vector<heap*> V;
  void build(){
     nullNd = new heap;
     nullNd->dep = 0;
     nullNd->edge = new nd;
     fill(nullNd->chd, nullNd->chd+4, nullNd);
     while(not dfsQ.empty()){
        int u = dfsQ.front(); dfsQ.pop();
        if(!nxt[ u ]) head[ u ] = nullNd;
else head[ u ] = head[nxt[ u ]->v];
        V.clear();
        for( auto&& e : g[ u ] ){
          int v = e \rightarrow v;
          if( dst[ v ] == -1 ) continue;
          e->d += dst[ v ] - dst[ u ];
if( nxt[ u ] != e ){
             heap* p = new heap;
             fill(p->chd, p->chd+4, nullNd);
             p->dep = 1;
             p->edge = e:
             .
V.push_back(p);
        if(V.empty()) continue;
       make_heap(V.begin(), V.end(), cmp);
#define L(X) ((\dot{X} << 1)+1)
#define R(X) ((X<<1)+2)
        for( size_t i = 0 ; i < V.size() ; i ++ ){
  if(L(i) < V.size()) V[i]->chd[2] = V[L(i)];
          else V[i]->chd[2]=nullNd;
          if(R(i) < V.size()) V[i]->chd[3] = V[R(i)];
          else V[i]->chd[3]=nullNd;
        head[u] = merge(head[u], V.front());
  }
  vector<LL> ans;
  void first_K(){
     ans.clear();
     priority_queue<node> Q;
     if( dst[ s ] == -1 ) return;
ans.push_back( dst[ s ] );
     if( head[s] != nullNd )
     Q.push(node(head[s], dst[s]+head[s]->edge->d));
for( int _ = 1 ; _ < k and not Q.empty() ; _ ++ ){
  node p = Q.top(), q; Q.pop();
  ans.push_back( p.d );
  if(head[n H > dge > y] | L = pullNd)[
        if(head[ p.H->edge->v ] != nullNd){
          q.H = head[p.H->edge->v];
          q.d = p.d + q.H->edge->d;
          Q.push(q);
       for( int i = 0 ; i < 4 ; i ++ )
  if( p.H->chd[ i ] != nullNd ){
    q.H = p.H->chd[ i ];
             q.d = p.d - p.H->edge->d + p.H->chd[i]->
                  edge->d;
             Q.push( q );
          }
     }
  }
```

```
void solve(){
    dijkstra();
    build();
    first_K();
}
solver;
```

## 5.14 Chordal Graph

```
struct Chordal {
  static const int MXN = 100010;
  vector<int> E[MXN], V[MXN];
  int n,f[MXN],rk[MXN],order(MXN],stk[MXN],nsz[MXN];
  bool vis[MXN], isMaximalClique[MXN];
  void init(int _n) {
     n = _n;
     for(int i = 0; i <= n; ++i) {
    E[i].clear(), V[i].clear();</pre>
       f[i]=rk[i]=order[i]=vis[i]=0;
    }
  void addEdge(int x, int y) {
    E[x].push_back(y), E[y].push_back(x);
  void mcs() {
     for(int i = 1; i <= n; ++i) V[0].push_back(i);
for(int i = n, M = 0; i >= 1; --i) {
       for(;;) {
         while(V[M].size()&&vis[V[M].back()])
           V[M].pop_back();
         if(V[M].size()) break; else M--;
       auto x=V[M].back();order[i]=x;rk[x]=i;vis[x]=1;
for(auto y : E[x]) if(!vis[y])
         f[y]++, V[f[y]].push_back(y), M=max(M,f[y]);
    }
  bool isChordal() {
     for(int i = 0; i <= n; ++i) vis[i] = stk[i] = 0;
for(int i = n; i >= 1; --i) {
       int top = 0, cnt = 0, m = n+1;
for(auto x : E[order[i]]) if(rk[x] > i)
         stk[top++]=x, vis[x]=1, m = min(m, rk[x]);
       if(m==n+1) continue
       for(auto x : E[order[m]]) if(vis[x]) ++cnt;
       for(int j = 0; j < top; ++j) vis[stk[j]] = 0;
if(cnt + 1 != top) return 0;</pre>
    return 1;
  void getMaximalClique() {
     for(int i = n; i >= 1; --i) {
       int M = n+1, w = order[i], v = 0;
       nsz[w] = 0; isMaximalClique[w] = 1;
       for(auto x : E[w]) if(rk[x] > i) {
         nsz[w]++;
         if(rk[x] < M) M = rk[x], v = x;
       if(v)isMaximalClique[v]&=nsz[v]+1>nsz[w];
    }
  int getMaximumClique() {
     int res = 0;
     for(int i = 1; i \le n; ++i) res=max(res,f[i]+1);
     return res;
  int getMaximumIndependentSet() {
     for(int i = 0; i <= n; ++i) vis[i] = 0;</pre>
     int res = 0;
     for(int i = 1; i <= n; ++i) if(!vis[order[i]]) {</pre>
       res++, vis[order[i]] = 1;
       for(auto x : E[order[i]]) vis[x] = 1;
     return res;
};
```

## 5.15 Graph Method

| Manhattan MST | For each point, consider the points that surround it(8 | octants). Then, connect it with the closest point. For example, consider 45~90. For each point p, the closest point is  $min\{x+y \mid x-y>=p.x-p.y, x>=p.x\}$ . Finally, the answer is this new graphs(E=4N) MST.

# 6 String

#### 6.1 PalTree

```
const int MXN = 1000010;
struct PalT{
  int nxt[MXN][26],fail[MXN],len[MXN];
  int tot,lst,n,state[MXN],cnt[MXN],num[MXN];
  int diff[MXN],sfail[MXN],fac[MXN],dp[MXN];
  char s[MXN] = \{-1\};
  int newNode(int l,int f){
    len[tot]=l,fail[tot]=f,cnt[tot]=num[tot]=0;
    memset(nxt[tot],0,sizeof(nxt[tot]));
diff[tot]=(l>0?l-len[f]:0);
    sfail[tot]=(l>0&&diff[tot]==diff[f]?sfail[f]:f);
    return tot++;
  int getfail(int x){
    while(s[n-len[x]-1]!=s[n]) x=fail[x];
    return x;
  int getmin(int v){
    dp[v]=fac[n-len[sfail[v]]-diff[v]];
    if(diff[v]==diff[fail[v]])
        dp[v]=min(dp[v],dp[fail[v]]);
    return dp[v]+1;
  int push(){
    int c=s[n]-'a',np=getfail(lst);
    if(!(lst=nxt[np][c])){
      lst=newNode(len[np]+2,nxt[getfail(fail[np])][c]);
      nxt[np][c]=lst; num[lst]=num[fail[lst]]+1;
    fac[n]=n;
    for(int v=lst;len[v]>0;v=sfail[v])
        fac[n]=min(fac[n],getmin(v));
    return ++cnt[lst],lst;
  void init(const char *_s){
    tot=lst=n=0;
    newNode(0,1),newNode(-1,1);
for(;_s[n];) s[n+1]=_s[n],++n,state[n-1]=push();
    for(int i=tot-1;i>1;i--) cnt[fail[i]]+=cnt[i];
}palt;
```

## 6.2 SAIS

```
const int N = 300010;
struct SA{
#define REP(i,n) for ( int i=0; i<int(n); i++ )</pre>
#define REP1(i,a,b) for (int i=(a); i <= int(b); i++)
  bool _t[N*2];
  int _s[N*2], _sa[N*2], _c[N*2], x[N], _p[N], _q[N*2],
       hei[N], r[N];
  int operator [] (int i){ return _sa[i]; }
  void build(int *s, int n, int m){
  memcpy(_s, s, sizeof(int) * n);
    sais(_s, _sa, _p, _q, _t, _c, n, m);
    mkhei(n);
  void mkhei(int n){
    REP(i,n) r[_sa[i]] = i;
    hei[0] = 0;
    REP(\bar{i},n) if(r[i]) {
      int ans = i>0 ? max(hei[r[i-1]] - 1, 0) : 0;
      while(_s[i+ans] == _s[_sa[r[i]-1]+ans]) ans++;
      hei[r[i]] = ans;
    }
  void sais(int *s, int *sa, int *p, int *q, bool *t,
      int *c, int n, int z){
    bool uniq = t[n-1] = true, neq;
    int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
         lst = -1;
#define MS0(x,n) memset((x),0,n*sizeof(*(x)))
#define MAGIC(XD) MS0(sa, n); \
```

```
memcpy(x, c, sizeof(int) * z); \
memcpy(x + 1, c, sizeof(int) * (z - 1)); \
REP(i,n) if(sa[i] \& !t[sa[i]-1]) sa[x[s[sa[i]-1]]++] =
      sa[i]-1; \
memcpy(x, c, sizeof(int) * z); \
for(int i = n - 1; i >= 0; i--) if(sa[i] && t[sa[i]-1])
      sa[--x[s[sa[i]-1]]] = sa[i]-1;
    MS0(c, z);
    REP(i,n) uniq \&= ++c[s[i]] < 2;
    REP(i,z-1) c[i+1] += c[i];
    if (uniq) { REP(i,n) sa[--c[s[i]]] = i; return; }
    for(int i = n - 2; i >= 0; i--) t[i] = (s[i]==s[i +1] ? t[i+1] : s[i]<s[i+1]);
    MAGIC(\overline{REP1}(\overline{i},1,\overline{n-1}) \overline{if}(\overline{t}[\overline{i}] 8\& !t[i-1]) sa[--x[s[i]]
         ]]]=p[q[i]=nn++]=i);
    REP(i, n) if (sa[i] && t[sa[i]] && !t[sa[i]-1]) {
      neq=lst<0|lmemcmp(s+sa[i],s+lst,(p[q[sa[i]]+1]-sa]
           [i])*sizeof(int));
      ns[q[lst=sa[i]]]=nmxz+=neq;
    sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz
          + 1);
    MAGIC(for(int i = nn - 1; i \ge 0; i--) sa[--x[s[p[
         nsa[i]]]] = p[nsa[i]];
}sa;
int H[N], SA[N], RA[N];
void suffix_array(int* ip, int len) {
  // should padding a zero in the back
  // ip is int array, len is array length
  // ip[0..n-1] != 0, and ip[len] = 0
  ip[len++] = 0;
  sa.build(ip, len, 128);
  memcpy(H,sa.hei+1,len<<2);</pre>
  memcpy(SA,sa._sa+1,len<<2)</pre>
  for(int i=0; i<len; i++) RA[i] = sa.r[i]-1;</pre>
  // resulting height, sa array \in [0,len)
```

#### 6.3 SuffixAutomata

```
// any path start from root forms a substring of S
// occurrence of P : iff SAM can run on input word P
// number of different substring : ds[1]-1
// total length of all different substring
                                                     dsl[1]
// max/min length of state i : mx[i]/mx[mom[i]]+1
// assume a run on input word P end at state i:
// number of occurrences of P : cnt[i]
// first occurrence position of P : fp[i]-IPI+1
// all position of P : fp of "dfs from i through rmom"
const int MXM = 1000010;
struct SAM{
  int tot, root, lst, mom[MXM], mx[MXM]; //ind[MXM]
int nxt[MXM][33]; //cnt[MXM],ds[MXM],dsl[MXM],fp[MXM]
  // bool v[MXM]
  int newNode(){
    int res = ++tot;
    fill(nxt[res], nxt[res]+33, 0);
    mom[res] = mx[res] = 0; //cnt=ds=dsl=fp=v=0
    return res;
  void init(){
    tot = 0; root = newNode(); lst = root;
  void push(int c){
    int p = lst;
    int np = newNode(); //cnt[np]=1
mx[np] = mx[p]+1; //fp[np]=mx[np]-1
for(; p && nxt[p][c] == 0; p = mom[p])
      nxt[p][c] = np;
     if(p == 0) mom[np] = root;
    else{
       int q = nxt[p][c];
       if(mx[p]+1 == mx[q]) mom[np] = q;
       else{
         int nq = newNode(); //fp[nq]=fp[q]
         mx[nq] = mx[p]+1;
         for(int i = 0; i < 33; i++)
           nxt[nq][i] = nxt[q][i];
         mom[nq] = mom[q]; mom[q] = nq; mom[np] = nq;
         for(; p && nxt[p][c] == q; p = mom[p])
```

```
nxt[p][c] = nq;
      }
    lst = np;
  void calc(){
    calc(root); iota(ind,ind+tot,1);
    sort(ind,ind+tot,[&](int i,int j){return mx[i]<mx[j</pre>
         ];});
    for(int i=tot-1;i>=0;i--)
    cnt[mom[ind[i]]]+=cnt[ind[i]];
  void calc(int x){
    v[x]=ds[x]=1;dsl[x]=0; //rmom[mom[x]].push_back(x);
    for(int i=0;i<26;i++){</pre>
      if(nxt[x][i]){
         if(!v[nxt[x][i]]) calc(nxt[x][i]);
ds[x]+=ds[nxt[x][i]];
         dsl[x]+=ds[nxt[x][i]]+dsl[nxt[x][i]];
    }
  }
  void push(char *str){
    for(int i = 0; str[i]; i++)
      push(str[i]-'a');
} sam;
```

#### 6.4 **Aho-Corasick**

```
struct ACautomata{
  struct Node{
     int cnt;
     Node *go[26], *fail, *dic;
     Node (){
       cnt = 0; fail = 0; dic=0;
       memset(go,0,sizeof(go));
  }pool[1048576],*root;
  int nMem;
  Node* new_Node(){
    pool[nMem] = Node();
     return &pool[nMem++];
  void init() { nMem = 0; root = new_Node(); }
  void add(const string &str) { insert(root,str,0); }
  void insert(Node *cur, const string &str, int pos){
  for(int i=pos;i<str.size();i++){</pre>
       if(!cur->go[str[i]-'a'])
         cur->go[str[i]-'a'] = new_Node();
       cur=cur->go[str[i]-'a'];
    cur->cnt++;
  void make_fail(){
     queue<Node*> que;
     que.push(root);
     while (!que.empty()){
       Node* fr=que.front(); que.pop();
for (int i=0; i<26; i++){
         if (fr->go[i]){
           Node *ptr = fr->fail;
           while (ptr && !ptr->go[i]) ptr = ptr->fail;
           fr->go[i]->fail=ptr=(ptr?ptr->go[i]:root);
           fr->go[i]->dic=(ptr->cnt?ptr:ptr->dic);
           que.push(fr->go[i]);
```

## 6.5 Z Value

```
void z_value(const char *s,int len,int *z){
  z[0]=len;
  for(int i=1,l=0,r=0;i<len;i++){</pre>
    z[i]=i < r?(i-l+z[i-l] < z[l]?z[i-l]:r-i):0;
    while(i+z[i]<len&&s[i+z[i]]==s[z[i]]) ++z[i];
    if(i+z[i]>r) l=i,r=i+z[i];
}
```

#### 6.6 BWT

```
struct BurrowsWheeler{
#define SIGMA 26
#define BASE 'a'
  vector<int> v[ SIGMA ];
  void BWT(char* ori, char* res){
    // make ori -> ori + ori
    // then build suffix array
}

void iBWT(char* ori, char* res){
  for( int i = 0 ; i < SIGMA ; i ++ )
    v[ i ].clear();
  int len = strlen( ori );
  for( int i = 0 ; i < len ; i ++ )
    v[ ori[i] - BASE ].push_back( i );
  vector<int> a;
  for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )
    for( auto j : v[ i ] ){
      a.push_back( j );
      ori[ ptr ++ ] = BASE + i;
    }
  for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
      res[ i ] = ori[ a[ ptr ] ];
      ptr = a[ ptr ];
    }
    res[ len ] = 0;
}
} bwt;</pre>
```

## 6.7 ZValue Palindrome

## 6.8 Smallest Rotation

```
//rotate(begin(s),begin(s)+minRotation(s),end(s))
int minRotation(string s) {
  int a = 0, N = s.size(); s += s;
  rep(b,0,N) rep(k,0,N) {
    if(a+k == b || s[a+k] < s[b+k])
      {b += max(0, k-1); break;}
    if(s[a+k] > s[b+k]) {a = b; break;}
  } return a;
}
```

## 6.9 Cyclic LCS

```
#define L 0
#define LU 1
#define U 2
const int mov[3][2]=\{0,-1,-1,-1,-1,0\};
int al,bl;
char a[MAXL*2],b[MAXL*2]; // 0-indexed
int dp[MAXL*2][MAXL];
char pred[MAXL*2][MAXL];
inline int lcs_length(int r) {
  int i=r+al, j=bl, l=0;
  while(i>r) {
    char dir=pred[i][j];
    if(dir==LU) l++
    i+=mov[dir][0]; j+=mov[dir][1];
  return 1;
inline void reroot(int r) { // r = new base row
  int i=r,j=1;
  while(j<=bl&&pred[i][j]!=LU) j++;</pre>
  if(j>bl) return;
  pred[i][j]=L;
  while(i<2*al&&j<=bl) {</pre>
    if(pred[i+1][j]==U) {
      i++; pred[i][j]=L;
```

```
} else if(j<bl&&pred[i+1][j+1]==LU) {</pre>
       i++; j++; pred[i][j]=L;
    } else j++;
  }
int cyclic_lcs() {
  \ensuremath{/\!/} a, b, al, bl should be properly filled
  // note: a WILL be altered in process

    concatenated after itself

  char tmp[MAXL];
  if(al>bl)
    swap(al,bl); strcpy(tmp,a);
     strcpy(a,b); strcpy(b,tmp);
  strcpy(tmp,a); strcat(a,tmp);
  // basic lcs
  for(int i=0;i<=2*al;i++)
     dp[i][0]=0; pred[i][0]=U;
  for(int j=0;j<=bl;j++) {</pre>
    dp[0][j]=0; pred[0][j]=L;
  for(int i=1;i<=2*al;i++)</pre>
    for(int j=1; j<=bl; j++) {
   if(a[i-1]==b[j-1]) dp[i][j]=dp[i-1][j-1]+1;</pre>
       else dp[i][j]=max(dp[i-1][j],dp[i][j-1]);
if(dp[i][j-1]==dp[i][j]) pred[i][j]=L;
       else if(\bar{a}[i-1]=b[j-1]) pred[i][j]=LU;
       else pred[i][j]=U;
  }
// do cyclic lcs
  int clcs=0;
  for(int i=0;i<al;i++) {</pre>
    clcs=max(clcs,lcs_length(i)); reroot(i+1);
  // recover a
  a[al]='\0':
  return clcs;
```

#### 7 Data Structure

## 7.1 Link-Cut Tree

```
const int MEM = 100005;
struct Splay {
  static Splay nil, mem[MEM], *pmem;
  Splay *ch[2], *f;
int val, rev, size;
  Splay (int _val=-1) : val(_val), rev(0), size(1)
  \{ f = ch[0] = ch[1] = &nil; \}
  bool isr()
  { return f->ch[0] != this && f->ch[1] != this; }
  int dir(){return f->ch[0] != this;}
void setCh(Splay *c, int d){
     ch[d] = c; if (c != &nil) c -> f = this; pull();
  void push(){
     if( !rev ) return;
    swap(ch[0], ch[1]);
if (ch[0] != &nil) ch[0]->rev ^= 1;
if (ch[1] != &nil) ch[1]->rev ^= 1;
     rev=0;
  void pull(){
     size = ch[0] - size + ch[1] - size + 1;
     if (ch[0] != &nil) ch[0]->f = this;
if (ch[1] != &nil) ch[1]->f = this;
}Splay::nil,Splay::mem[MEM],*Splay::pmem=Splay::mem;
Splay *nil = &Splay::nil;
void rotate(Splay *x){
  Splay *p = x \rightarrow f; int d = x \rightarrow dir();
  if (!p->isr()) p->f->setCh(x, p->dir());
  else x->f = p->f
  p->setCh(x->ch[!d], d); x->setCh(p, !d);
vector<Splay*> splayVec;
void splay(Splay *x){
  splayVec.clear();
  for (Splay *q=x;; q=q->f){
```

```
splayVec.push_back(q);
    if (q->isr()) break;
 reverse(begin(splayVec), end(splayVec));
for (auto it : splayVec) it->push();
while (!x->isr()) {
    if (x->f->isr()) rotate(x);
    else if (x->dir()==x->f->dir())
      rotate(x->f),rotate(x);
    else rotate(x),rotate(x);
int id(Splay *x) { return x - Splay::mem + 1; }
Splay* access(Splay *x){
  Splay *q = nil;
  for (;x!=nil;x=x->f){
    splay(x); x-setCh(q, 1); q = x;
  return q;
}
void chroot(Splay *x){
 access(x); splay(x); x->rev ^= 1;
void link(Splay *x, Splay *y){
  chroot(y); y->f=x;
void cut_p(Splay *y) {
 access(y); splay(y); y->ch[0] = y->ch[0]->f = nil;
void cut(Splay *x, Splay *y){
  chroot(x); cut_p(y);
Splay* get_root(Splay *x) {
  x=access(x);
  for(; x \rightarrow ch[0] != nil; x = x \rightarrow ch[0]) x \rightarrow push();
  splay(x); return x;
bool conn(Splay *x, Splay *y) {
  return get_root(x) == get_root(y);
Splay* lca(Splay *x, Splay *y) {
  access(x); return access(y);
/* query(Splay *x,Splay *y){
  setroot(y),x=access(x); return x->size;
/* query(Splay *x,Splay *y){
  Splay *p=lca(x,y);
  return p \rightarrow val + p \rightarrow ch[1] \rightarrow size + (x! = p?x \rightarrow size : 0);
```

# 8 Others

## 8.1 Find max tangent(x,y is increasing)

```
const int MAXN = 100010;
Pt sum[MAXN], pnt[MAXN], ans, calc;
inline bool cross(Pt a, Pt b, Pt c){
  return (c.y-a.y)*(c.x-b.x) > (c.x-a.x)*(c.y-b.y);
}//pt[0]=(0,0);pt[i]=(i,pt[i-1].y+dy[i-1]),i=1~n;dx>=l
double find_max_tan(int n,int l,LL dy[]){
  int np, st, ed, now;
  sum[0].x = sum[0].y = np = st = ed = 0;
for (int i = 1, v; i <= n; i++)
    sum[i].x=i,sum[i].y=sum[i-1].y+dy[i-1];
  ans.x = now = 1,ans.y = -1;
for (int i = 0; i <= n - 1; i++){
     while(np>1&&cross(pnt[np-2],pnt[np-1],sum[i]))
     if (np < now \&\& np != 0) now = np;
    pnt[np++] = sum[i];
     while(now<np&!cross(pnt[now-1],pnt[now],sum[i+l]))</pre>
     calc = sum[i + l] - pnt[now - 1];
     if (ans.y * calc.x < ans.x * calc.y)
       ans = calc,st = pnt[now - 1].x,ed = i + l;
  return (double)(sum[ed].y-sum[st].y)/(sum[ed].x-sum[
       st].x);
}
```

## 8.2 Exact Cover Set

```
// given n*m 0-1 matrix
// find a set of rows s.t.
// for each column, there's exactly one 1
#define N 1024 //row
#define M 1024 //column
#define NM ((N+2)*(M+2))
char A[N][M]; //n*m 0-1 matrix
bool used[N]; //answer: the row used
int id[N][M]
int L[NM],R[NM],D[NM],U[NM],C[NM],S[NM],ROW[NM];
void remove(int c){
  L[R[c]]=L[c]; R[L[c]]=R[c]
  for( int i=D[c]; i!=c; i=D[i] )
  for( int j=R[i]; j!=i; j=R[j] )
       U[D[j]]=U[j]; D[U[j]]=D[j]; S[C[j]]--;
void resume(int c){
  for( int i=D[c]; i!=c; i=D[i] )
  for( int j=L[i]; j!=i; j=L[j] ){
       U[D[j]]=D[U[j]]=j; S[C[j]]++;
  L[R[c]]=R[L[c]]=c;
bool dfs(){
  if(R[0]==0) return 1;
  int md=100000000,c;
  for( int i=R[0]; i!=0; i=R[i] )
  if(S[i] < md) { md=S[i]; c=i; }</pre>
  if(md==0) return 0;
  remove(c);
  for( int i=D[c]; i!=c; i=D[i] ){
     used[ROW[i]]=1;
     for( int j=R[i]; j!=i; j=R[j] ) remove(C[j]);
if(dfs()) return 1;
     for( int j=L[i]; j!=i; j=L[j] ) resume(C[j]);
     used[ROW[i]]=0;
  resume(c);
  return 0;
bool exact_cover(int n,int m){
  for( int i=0; i<=m; i++ ){
   R[i]=i+1; L[i]=i-1; U[i]=D[i]=i;</pre>
     S[i]=0; C[i]=i;
  R[m]=0; L[0]=m;
  int t=m+1;
  for( int i=0; i<n; i++ ){</pre>
     int k=-1;
for( int j=0; j<m; j++ ){
       if(!A[i][j]) continue;
if(k=-1) L[t]=R[t]=t;
else{ L[t]=k; R[t]=R[k];
        k=t; D[t]=j+1; U[t]=U[j+1];
        L[R[t]]=R[L[t]]=U[D[t]]=D[U[t]]=t;
        C[t]=j+1; S[C[t]]++; ROW[t]=i; id[i][j]=t++;
  for( int i=0; i<n; i++ ) used[i]=0;</pre>
  return dfs();
}
```

## 8.3 Binary Next Permutation

```
ull next_perm(ull v){
  ull t=v|(v-1);
  return (t+1)|(((~t&-~t)-1)>>(__builtin_ctzll(v)+1));
}
```

#### 8.4 Hilbert Curve

```
long long hilbert(int n, int x, int y) {
   long long res = 0;
   for (int s = n / 2; s; s >>= 1) {
      int rx = (x & s) > 0;
      int ry = (y & s) > 0;
      res += s * 111 * s * ((3 * rx) ^ ry);
      if (ry == 0) {
        if (rx == 1) x = s - 1 - x, y = s - 1 - y;
      }
}
```

```
swap(x, y);
}
return res;
}
```