

CHAPTER 1

1.1 Base-10: 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32
 Octal: 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 40
 Hex: 10 11 12 13 14 15 16 17 18 19 1A 1B 1C 1D 1E 1F 20
 Base-13 A B C 10 11 12 13 14 15 16 17 18 19 23 24 25 26

1.2 (a) 32,768 (b) 67,108,864 (c) 6,871,947,674

1.3 $(4310)_5 = 4 * 5^3 + 3 * 5^2 + 1 * 5^1 = 580_{10}$
 $(198)_{12} = 1 * 12^2 + 9 * 12^1 + 8 * 12^0 = 260_{10}$
 $(735)_8 = 7 * 8^2 + 3 * 8^1 + 5 * 8^0 = 477_{10}$
 $(525)_6 = 5 * 6^2 + 2 * 6^1 + 5 * 6^0 = 197_{10}$

1.4 14-bit binary: 11_1111_1111_1111
 Decimal: $2^{14} - 1 = 16,383_{10}$
 Hexadecimal: 3FFF₁₆

1.5 Let b = base

(a) $14/2 = (b + 4)/2 = 5$, so b = 6

(b) $54/4 = (5*b + 4)/4 = b + 3$, so $5 * b = 52 - 4$, and b = 8

(c) $(2 * b + 4) + (b + 7) = 4b$, so b = 11

1.6 $(x - 3)(x - 6) = x^2 - (6 + 3)x + 6*3 = x^2 - 11x + 22$

Therefore: $6 + 3 = b + 1$ so b = 8

Also, $6*3 = (18)_{10} = (22)_8$

1.7 $68BE = 0110_1000_1011_1110 = 110_100_010_111_110 = (64276)_8$

1.8 (a) Results of repeated division by 2 (quotients are followed by remainders):

$431_{10} = 215(1); 107(1); 53(1); 26(1); 13(0); 6(1) 3(0) 1(1)$

Answer: $1111_1010_2 = FA_{16}$

(b) Results of repeated division by 16:

$431_{10} = 26(15); 1(10)$ (Faster)

Answer: $FA = 1111_1010$

1.9 (a) $10110.0101_2 = 16 + 4 + 2 + .25 + .0625 = 22.3125$

(b) $16.5_{16} = 16 + 6 + 5*(.0615) = 22.3125$

(c) $26.24_8 = 2 * 8 + 6 + 2/8 + 4/64 = 22.3125$

(d) $FAFA.B_{16} = 15 \cdot 16^3 + 10 \cdot 16^2 + 15 \cdot 16 + 10 + 11/16 = 64,250.6875$

(e) $1010.1010_2 = 8 + 2 + .5 + .125 = 10.625$

1.10 (a) $1.10010_2 = 0001.1001_2 = 1.9_{16} = 1 + 9/16 = 1.563_{10}$

(b) $110.010_2 = 0110.0100_2 = 6.4_{16} = 6 + 4/16 = 6.25_{10}$

Reason: 110.010_2 is the same as 1.10010_2 shifted to the left by two places.

1.11

$$\begin{array}{r}
 1011.11 \\
 101 \overline{) 111011.0000} \\
 \underline{101} \\
 01001 \\
 \underline{101} \\
 1001 \\
 \underline{101} \\
 1000 \\
 \underline{101} \\
 0110
 \end{array}$$

The quotient is carried to two decimal places, giving 1011.11

Checking: $111011_2 / 101_2 = 59_{10} / 5_{10} \cong 1011.11_2 = 58.75_{10}$

1.12 (a) 10000 and 110111

$$\begin{array}{r}
 1011 \\
 +101 \\
 \hline
 10000 = 16_{10}
 \end{array}$$

$$\begin{array}{r}
 1011 \\
 \times 101 \\
 \hline
 1011 \\
 1011 \\
 \hline
 110111 = 55_{10}
 \end{array}$$

(b) 62_h and 958_h

$$\begin{array}{r}
 2E_h \quad 0010_1110 \\
 +34_h \quad 0011_0100 \\
 \hline
 62_h \quad 0110_0010 = 98_{10}
 \end{array}$$

$$\begin{array}{r}
 2E_h \\
 \times 34_h \\
 \hline
 B^38 \\
 \hline
 8^2A \\
 \hline
 9 \ 5 \ 8_h = 2392_{10}
 \end{array}$$

1.13 (a) Convert 27.315 to binary:

	Integer Quotient		Remainder	Coefficient
27/2 =	13	+	$\frac{1}{2}$	$a_0 = 1$
13/2	6	+	$\frac{1}{2}$	$a_1 = 1$
6/2	3	+	0	$a_2 = 0$
3/2	1	+	$\frac{1}{2}$	$a_3 = 1$
$\frac{1}{2}$	0	+	$\frac{1}{2}$	$a_4 = 1$

$$27_{10} = 11011_2$$

	Integer	Fraction	Coefficient
.315 x 2	= 0	+ .630	$a_{-1} = 0$
.630 x 2	= 1	+ .26	$a_{-2} = 1$
.26 x 2	= 0	+ .52	$a_{-3} = 0$
.52 x 2	= 1	+ .04	$a_{-4} = 1$

$$.315_{10} \cong .0101_2 = .25 + .0625 = .3125$$

$$27.315 \cong 11011.0101_2$$

(b) $2/3 \cong .666666667$

	Integer	Fraction	Coefficient
.6666_6666_67 x 2	= 1	+ .3333_3333_34	$a_{-1} = 1$
.333333334 x 2	= 0	+ .666666668	$a_{-2} = 0$
.666666668 x 2	= 1	+ .333333336	$a_{-3} = 1$
.333333336 x 2	= 0	+ .666666672	$a_{-4} = 0$
.666666672 x 2	= 1	+ .333333344	$a_{-5} = 1$
.333333344 x 2	= 0	+ .666666688	$a_{-6} = 0$
.666666688 x 2	= 1	+ .333333376	$a_{-7} = 1$
.333333376 x 2	= 0	+ .666666752	$a_{-8} = 0$

$$.666666667_{10} \cong .10101010_2 = .5 + .125 + .0313 + .0078 = .6641_{10}$$

$$.10101010_2 = .1010_2 = .AA_{16} = 10/16 + 10/256 = .6641_{10} \text{ (Same as (b)).}$$

1.14	(a)	1000_0000 1s comp: 0111_1111 2s comp: 1000_0000	(b)	0000_0000 1s comp: 1111_1111 2s comp: 0000_0000	(c)	1101_1010 1s comp: 0010_0101 2s comp: 0010_0110
	(d)	0111_0110 1s comp: 1000_1001 2s comp: 1000_1010	(e)	1000_0101 1s comp: 0111_1010 2s comp: 0111_1011	(f)	1111_1111 1s comp: 0000_0000 2s comp: 0000_0001

1.15	(a)	52,784,630 9s comp: 47,215,369 10s comp: 47,215,370	(b)	63,325,600 9s comp: 36,674,399 10s comp: 36,674,400
	(c)	25,000,000 9s comp: 74,999,999 10s comp: 75,000,000	(d)	00,000,000 9s comp: 99,999,999 10s comp: 00,000,000

1.16	B2FA	B2FA: 1011_0010_1111_1010
	15s comp: 4D05	1s comp: 0100_1101_0000_0101
	16s comp: 4D06	2s comp: 0100_1101_0000_0110 = 4D06

1.17 **(a)** $3409 \rightarrow 03409 \rightarrow 96590 \text{ (9s comp)} \rightarrow 96591 \text{ (10s comp)}$
 $06428 - 03409 = 06428 + 96591 = 03019$

(b) $1800 \rightarrow 01800 \rightarrow 98199 \text{ (9s comp)} \rightarrow 98200 \text{ (10 comp)}$
 $125 - 1800 = 00125 + 98200 = 98325 \text{ (negative)}$
 Magnitude: 1675
 Result: $125 - 1800 = 1675$

(c) $6152 \rightarrow 06152 \rightarrow 93847$ (9s comp) $\rightarrow 93848$ (10s comp)
 $2043 - 6152 = 02043 + 93848 = 95891$ (Negative)
 Magnitude: 4109
 Result: $2043 - 6152 = -4109$

(d) $745 \rightarrow 00745 \rightarrow 99254$ (9s comp) $\rightarrow 99255$ (10s comp)
 $1631 - 745 = 01631 + 99255 = 0886$ (Positive)
 Result: $1631 - 745 = 886$

1.18 Note: Consider sign extension with 2s complement arithmetic.

<p>(a)</p> $\begin{array}{r} 10001 \\ 1s \text{ comp: } 01110 \\ 2s \text{ comp: } 01111 \\ \hline 10011 \\ \text{Diff: } 00010 \end{array}$	<p>(b)</p> $\begin{array}{r} 100011 \\ 1s \text{ comp: } 1011100 \text{ with sign extension} \\ 2s \text{ comp: } 1011101 \\ \hline 0100010 \\ 1111111 \text{ sign bit indicates that the result is negative} \\ 0000001 \text{ 2s complement} \\ -000001 \text{ result} \end{array}$
<p>(c)</p> $\begin{array}{r} 101000 \\ 1s \text{ comp: } 1010111 \\ 2s \text{ comp: } 1011000 \\ \hline 001001 \\ \text{Diff: } 1100001 \text{ (negative)} \\ 0011111 \text{ (2s comp)} \\ -011111 \text{ (diff is -31)} \end{array}$	<p>(d)</p> $\begin{array}{r} 10101 \\ 1s \text{ comp: } 1101010 \text{ with sign extension} \\ 2s \text{ comp: } 1101011 \\ \hline 110000 \\ 0011011 \text{ sign bit indicates that the result is positive} \\ \text{Check: } 48 - 21 = 27 \end{array}$

1.19 $+9286 \rightarrow 009286$; $+801 \rightarrow 000801$; $-9286 \rightarrow 990714$; $-801 \rightarrow 999199$

(a) $(+9286) + (-801) = 009286 + 000801 = 010087$

(b) $(+9286) + (-801) = 009286 + 999199 = 008485$

(c) $(-9286) + (+801) = 990714 + 000801 = 991515$

(d) $(-9286) + (-801) = 990714 + 999199 = 989913$

1.20 $+49 \rightarrow 0_110001$ (Needs leading zero indicate + value); $+29 \rightarrow 0_011101$ (Leading 0 indicates + value)
 $-49 \rightarrow 1_001111$; $-29 \rightarrow 1_100011$

(a) $(+29) + (-49) = 0_011101 + 1_001111 = 1_101100$ (1 indicates negative value.)
 Magnitude = 0_010100 ; Result $(+29) + (-49) = -20$

(b) $(-29) + (+49) = 1_100011 + 0_110001 = 0_010100$ (0 indicates positive value)
 $(-29) + (+49) = +20$

(c) Must increase word size by 1 (sign extension) to accomodate overflow of values:
 $(-29) + (-49) = 11_100011 + 11_001111 = 10_110010$ (1 indicates negative result)
 Magnitude: $1_001110 = 78_{10}$
 Result: $(-29) + (-49) = -78$

1.21 +9742 → 009742 → 990257 (9's comp) → 990258 (10s comp)
+641 → 000641 → 999358 (9's comp) → 999359 (10s comp)

(a) (+9742) + (+641) → 010383

(b) (+9742) + (-641) → 009742 + 999359 = 009102
Result: (+9742) + (-641) = 9102

(c) -9742 + (+641) = 990258 + 000641 = 990899 (negative)
Magnitude: 009101
Result: (-9742) + (+641) = -9101

(d) (-9742) + (-641) = 990258 + 999359 = 989617 (Negative)
Magnitude: 10383
Result: (-9742) + (-641) = -10383

1.22 8,723
BCD: 1000_0111_0010_0011
ASCII: 0_011_1000_011_0111_011_0010_011_0001

1.23

1000	0100	0010	(842)
<u>0101</u>	<u>0011</u>	<u>0111</u>	(+537)
1101	0111	1001	
<u>0110</u>			
0001	0011	0111	0101 (1,379)

(a)	(b)
6 3 1 1 Decimal	6 4 2 1 Decimal
0 0 0 0 0	0 0 0 0 0
0 0 0 1 1	0 0 0 1 1
0 0 1 0 2	0 0 1 0 2
0 1 0 0 3	0 0 1 1 3
0 1 1 0 4 (or 0101)	0 1 0 0 4
0 1 1 1 5	0 1 0 1 5
1 0 0 0 6	1 0 0 0 6 (or 0110)
1 0 1 0 7 (or 1001)	1 0 0 1 7
1 0 1 1 8	1 0 1 0 8
1 1 0 0 9	1 0 1 1 9

1.25

(a) 5,137 ₁₀	BCD: 0101_0011_0111
(b)	Excess-3: 1000_0100_0110_1010
(c)	2421: 1011_0001_0011_0111
(d)	6311: 0111_0001_0100_1001

1.26 5,137 9s Comp: 4,862
2421 code: 0100_1110_1100_1000
1s comp: 1011_0001_0011_0111 same as (c) in 1.25

1.27 For a deck with 52 cards, we need 6 bits ($32 < 52 < 64$). Let the msb's select the suit (e.g., diamonds, hearts, clubs, spades are encoded respectively as 00, 01, 10, and 11. The remaining four bits select the "number" of the card. Example: 0001 (ace) through 1011 (9), plus 101 through 1100 (jack, queen, king). This a jack of spades might be coded as 11_1010. (Note: only 52 out of 64 patterns are used.)

1.28 G (dot) (space) B o o l e
01000111_11101111_01101000_01101110_00100000_11000100_11101111_11100101

1.29 Bill Gates

1.30 73 F4 E5 76 E5 4A EF 62 73

73: 0_111_0011 s
F4: 1_111_0100 t
E5: 1_110_0101 e
76: 0_111_0110 v
E5: 1_110_0101 e
4A: 0_100_1010 j
EF: 1_110_1111 o
62: 0_110_0010 b
73: 0_111_0011 s

1.31 $62 + 32 = 94$ printing characters

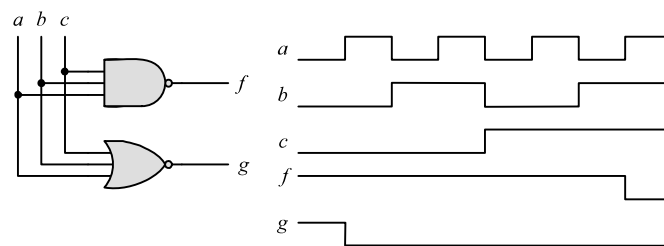
1.32 bit 6 from the right

1.33 (a) 897 (b) 564 (c) 871 (d) 2,199

1.34 ASCII for decimal digits with odd parity:

(0): 10110000 (1): 00110001 (2): 00110010 (3): 10110011
(4): 00110100 (5): 10110101 (6): 10110110 (7): 00110111
(8): 00111000 (9): 10111001

1.35 (a)



1.36

