CS101 Introduction to computing

Recursive Function and Problem solving using function

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Outline

- Recursive function
- Mutual recursion
- Problem solving using recursions
- Recursion: using Array

Fib (N): Number of Recursive Call

Multiple Recursive Calls

$$Fib(N) = Fib(n-1) + Fib(n-2)$$

- Number of recursive call for N
 - Claim: Number of recursive call for Fib(n-1) is higher than number of recursive call for Fib(n-2)
 - Denote number of recursive call for Fib(n) = f_n
 - $-f_{n-1} > f_{n-2}$

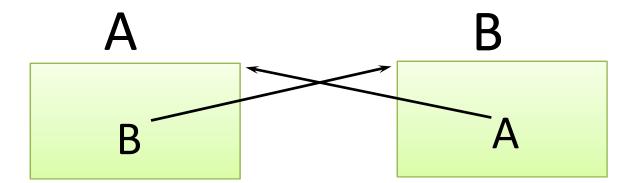
Fib (N): Number of Recursive Call

- Can I Say: $f_n = f_{n-1} + f_{n-2} > f_{n-2} + f_{n-2} = 2. f_{n-2}$
- Then $f_n > 2.f_{n-2} > 2.2.f_{n-4} > 2.2.2.f_{n-6}$ =...= $2^{n/2} f_1$ So $f_n > 2^{n/2}$
- Number of recursive call require to compute Fib(n) is > 2^{n/2}
- Can you calculate for 200 Fibonacci using recursive program
 - Will take at least 2¹⁰⁰ recursive call: huge time and space

Mutual Recursion

Recursion doesn't always occur because a routine calls itself...

Mutual Recursion occurs when two routines call each other.



Mutual Recursion Example EvenOdd

- Problem: Determine whether a number,
 N, is odd or even.
 - —If N is equal to 0, then n is even
 - —N is odd if N-1 is even

```
int Even(int N){
  if(n==0) return 1;
  return Odd(n-1);
}
```

```
int Odd(int N){
  if(n==0) return 0;
  return Even(n-1);
}
```

Recursion Example: Problem Solving

```
int Reverse(int n) {
int RevNum=0, Rem;
RevNum=0;
while(n != 0)
  Rem = n%10;
 RevNum=RevNum*10+Rem;
 n=n/10;
 return RevNum;
```

```
int Reverse(int n) {
static int Rev=0;
if(n ==0) return 0;
Rev = Rev *10;
Rev = Rev + N%10;
Reverse(N/10);
return Rev;
```

<u>GCD</u>

- By definition from Euclid's algorithm
- Recursive

```
-GCD(a,b) if a > b GCD(a%b, b)
if b > a GCD(a, b%a)
```

Code looks simpler

Recursive GCD

Code looks simpler

```
int GCD(int a, int b) {
  if(a==0) return b;
  return gcd(b%a, a);
}
```

- Trace 35, 10 == > 10, 35 == > 5, 10 == > 0, 5
- Trace 10, 15 == > 5, 10 == > 0, 5
- Trace 31, 2 == > 2, 31 == > 1, 2 == > 0, 1

Modular C Code: Binary Search

```
int BinSrch
 (int Rmin,
  int Rmax, int X){
while (Rmin<Rmax) {</pre>
mid=(Rmin+Rmax)/2;
  if(X==mid)
    return mid;
  if (X>mid)
      Rmin=mid+1;
  else Rmax=mid;
  return -1;
```

```
int BinSrch(int Rmin,
     int Rmax, int X){
int mid;
mid=(Rmin+Rmax)/2;
if(X==mid)return mid;
if (X>mid)
return
  BinSrch(mid+1,Rmax,X);
else
return BinSrch(Rmin, mid, X);
  return -1;
```

Binary Search Analysis

```
• BinSrch(Range, X)
 - mid=Range/2;
 - BinSrch(Range/2,X);
\bullet B(R) = B(R/2) + c
         = B(R/4) + 2c
         = B(1) + log_2R.c
         = c.ceil(log_2R)
```

Problem Solving Array, Function, Recursion

Problem Solving: Examples

- Max of an Array
 - -Iterative and recursive procedure
- Sieve of Eratos-thenes
- Array Reversal
- Sorting an array
 - -Bubble sort

Maximum of an Array

Max of an Array

Iterative Approach

```
int MaxOfAnArray(int A[], int n){
  int i, L=A[0];
  for(i=1; i<n; i++)
    if(A[i]>L)
       L=A[i];
  return L;
}
```

- Number of steps required
 - N: linear code

Max of an Array: Recursive Approach

```
Max( 2, 5, 12, 8, 16, 23, 1, 5)

= Max (Max of 1<sup>st</sup> Half, Max of 2<sup>nd</sup> Half)

= Max (Max(2,5,12,8), Max(16,23,1,5))
```

```
int RMax(int A[], int n1, int n2){
 int L1, L2;
  if(n1==n2) return A[n1];
 L1=Rmax(A, n1, (n1+n2)/2);
 L2=Rmax(A,(n1+n2)/2, n2);
  if (L1>L2) return L1;
  else return L2;
```

Max of an Array: Recursive Approach

Number of steps required

```
-R(N) = R(N/2) + R(N/2) + C
= 2.R(N/2) + C
= 4.R(N/4) + 2C = 8.R(N/8) + 4C = N.R(1) + N/2.C
= N + N/2.C = linear number of steps
```

```
int RMax(int A[], int n1, int n2){
int L1, L2; if(n1==n2) return A[n1];
L1=Rmax(A, n1, (n1+n2)/2);
L2=Rmax(A,(n1+n2)/2, n2);
if (L1>L2) return L1; else return L2;
}
```

Sieve of Eratosthenes

What is the sieve of Eratosthenes?

- Used to find prime number between 2 and N
- It works by gradually eliminating multiple of smallest unmark prime (x) in the given interval [2-N]
 - $-Till x^2 > N$
- Let us see with an example

Sieve of Eratosthenes: 1 to 20

Current Prime 2

[2,3,4,5,6,7,8,9,10,11,12,13, 14,15,16,17,18,19,20]

Current Prime 2

[2,3,4,5,6,7,8,9,10,11,12,13,14,15,16,17,18,19,20]

Current Prime 3

List: [2, 3, 5, 7, 9, 11, 13, 15, 17, 19]

• Current Prime 3

List: [2, 3, 5, 7, 9, 11,13,15, 17,19]

Current Prime 5 and 5< ceil(sqrt(20))

List: [2, 3, 5, 7, 11, 13, 17,19] //All are primes

Sieve of Eratosthenes

```
#define STRIKED 0 //Happened to be composite
#define NONSTRIKED 1 //Assumed Prime
void SeiveOfEratosthenes(int prime[MaxNum],
                  int N){
 int i,j, CP, SqrtM=sqrt(MaxNum)+1;
 for(i=0;i<MaxNum;i++) prime[i] = NONSTRIKED;</pre>
 for(i=2;i<SqrtM;i++){</pre>
   if (prime[i] == STRIKED) continue ;
   CP=i; //current Prime
   for(j=2*CP;j<MaxNum; j=j+CP) //do striking</pre>
         prime[j] = STRIKED;
```

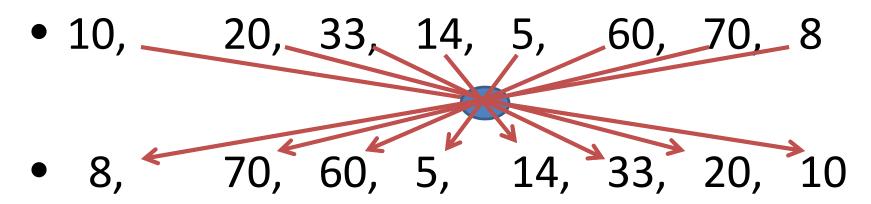
Reversing an Array

Reversing an Array

Approach

- Input: 10, 20, 33, 14, 5, 60, 70, 8
- Output: 8, 70, 60, 5, 14, 33, 20, 10
- Mirror Image
- 10, 20, 33, 14, 5, 60, 70, 8 | 8, 70, 60, 5, 14, 33, 20, 10
- 10, 20, 33, 14, 5, 60, 70, 8
- 8, 70, 60, 5, 14, 33, 20, 10

Reversing an Array



- Start from both ends: Exchange element
- Repeat for next elements form both sides till they meet each other

Reversing an Array: Iterative Code

```
void Reverse(int A[], int N,){
   int tmp,i=0,j=N-1;
   while(i<j){
      tmp=A[i]; [i]=A[j];A[j]=tmp;
      i=i+1,
      j=j-1;
   }
}</pre>
```

Reversing an Array: Recursive Code

```
void Reverse(int A[],
            int i, int j){
int tmp;
if(i<j){
    tmp=A[i]; A[i]=A[j];A[j]=tmp;
    Reverse(A, i+1, j-1);
void ReverseArray(int A[], int N){
    Reverse(A, 0, N-1);
```

Sorting an Array

Sorting an Array

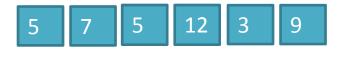
- Sorting: Arrange the element in some specific order
- Simpler case
 - non-decreasing : smallest element at beginning and biggest element at end
 - non-increasing order: biggest element at beginning and smallest element at end
- Example Non-sorted

30, 4, 7, 10, 12, 8, 2, 8

• Sorted : non-decreasing

2, 4, 7, 8, 8, 10, 12, 30

- 1. Start at the beginning of the data set.
- 2. Compare the first two elements, and if the first is greater than the second, swaps them.
- 3. Continue doing this for each pair of adjacent elements to the end of the data set.
- 4. Start again with the first two elements, repeating until no swaps have occurred on the last pass.



Invariant

Last PassNumber elements will be in sorted order

Last 0 elements are in sorted order



Pass 2 5 5 7 3 9 12 5 5 7 3 9 12 5 5 7 3 9 12 5 5 7 3 9 12

Invariant

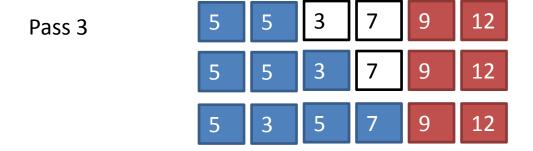
Last PassNumber elements will be in sorted order

Last 1 element is in sorted order
Show in RED Box



Invariant

Last PassNumber elements will be in sorted order



Last 2 elements are in sorted order
Show in RED Box



Pass 4 5 3 5 7 9 12 3 5 7 9 12

Invariant

Last PassNumber elements will be in sorted order

Last 3 elements are in sorted order
Show in RED Box



Pass 5 3 5 5 7 9 12 3 5 5 7 9 12

Invariant

Last PassNumber elements will be in sorted order

Last 4 elements are in sorted order
Show in RED Box

Sorted !!!

12

After 3 5 5 7
Pass 5

Invariant

Last PassNumber elements will be in sorted order

All 5 elements are in sorted order
Show in RED Box

Sorted !!!

Bubble Sort: C Code

```
void BubbleSort(int A[], int N){
   int Pass, j, tmp;
   for(Pass=0; Pass<N; Pass++){</pre>
     for(j=0; j<(N-Pass-1); j++){
        if(A[j] > A[j+1])
              tmp=A[j];
              A[j] = A[j+1];
              A[j+1]=tmp;
   }//end Pass
 //end BubbleSort
```

Thanks

Good luck for your MidSem Exam

