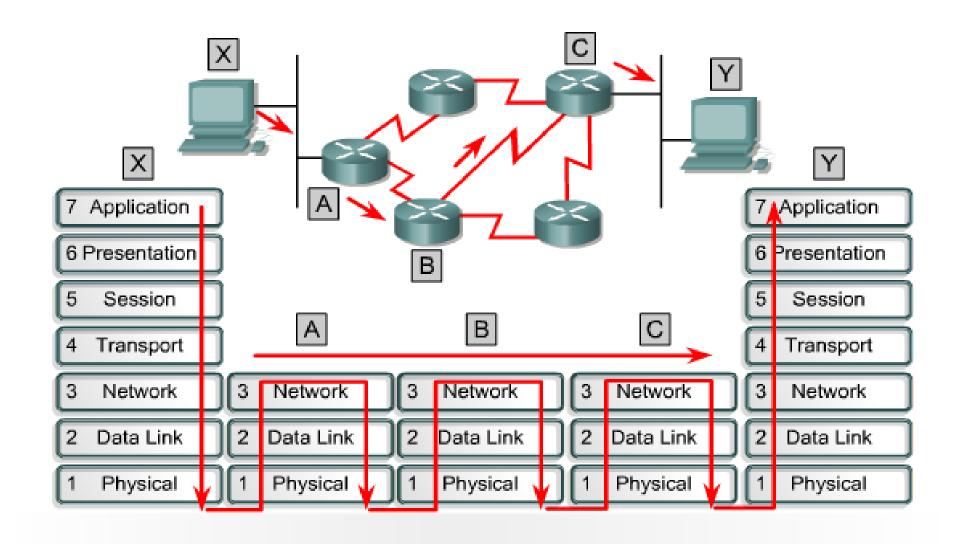
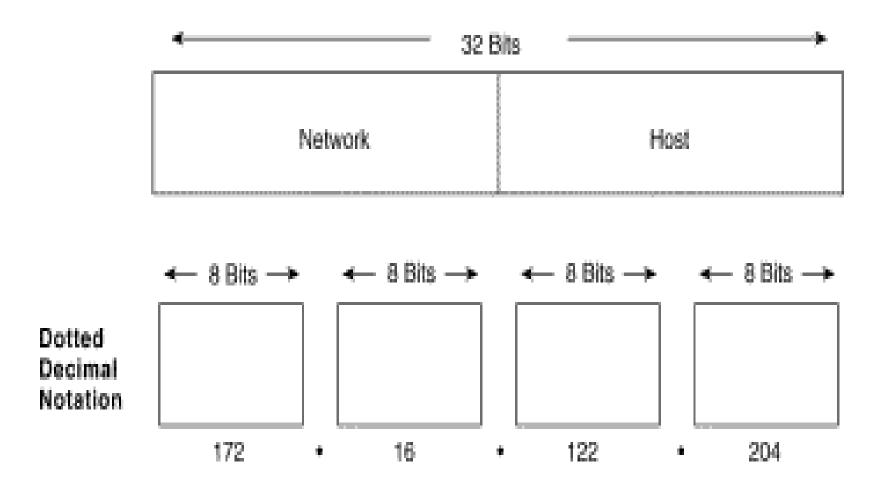
# IP Addressing

- IP is a connectionless, unreliable, best-effort delivery protocol.
- All nodes are identified using an IP address.
- IP accepts whatever data is passed down to it from the upper layers and forwards the data in the form of IP packets.
- Packets are delivered from the source to the destination using IP address.
- IP address is for the interface of a host.
- Multiple interfaces mean multiple IP addresses.



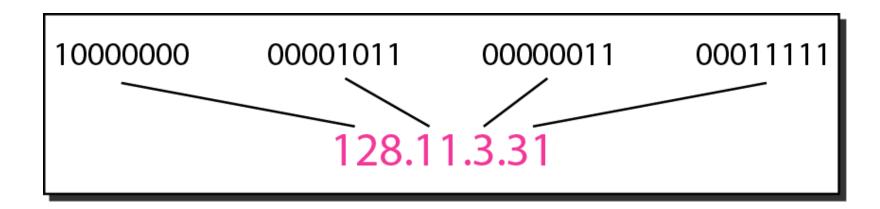
Each router provides its services to support upper-layer functions.

## **IP Address**



### IPv4 Addresses

- IPv4 address is a 32 bit address.
- IPv4 addresses are unique and universal (all nodes connecting Internet must have IP addresses).
- Address space of IPv4 is 2<sup>32</sup>



# Classful Addressing

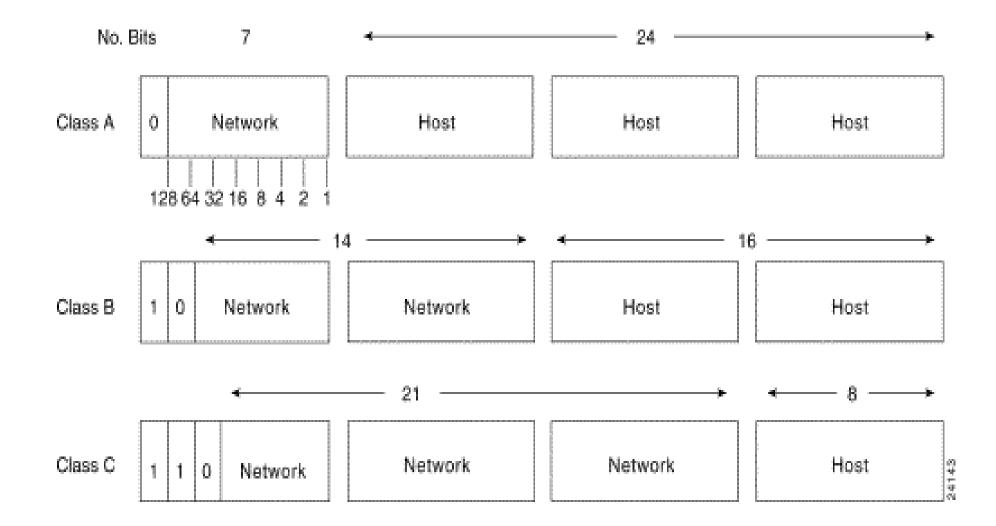
- The address space is divided into five classes: A, B, C, D, and E.
- A large part of the available addresses are wasted.
- Almost obsolete.

	First byte	Second byte	Third byte	Fourth byte
Class A	0			
Class B	10			
Class C	110			
Class D	1110			
Class E	1111			

	First byte	Second byte	Third byte	Fourth byte
Class A	0–127			
Class B	128–191			
Class C	192–223			
Class D	224–239			
Class E	240–255			

a. Binary notation

b. Dotted-decimal notation



### Number of blocks and block size in classful IPv4 addressing

Class	Number of Blocks	Block Size	Application
A	128	16,777,216	Unicast
В	16,384	65,536	Unicast
С	2,097,152	256	Unicast
D	1	268,435,456	Multicast
Е	1	268,435,456	Reserved

- Class A: Address begins with bit 0. It has 8 bit network number (range 0.0.0.0-to-127.255.255.255), 24 bit host number.
- Class B: Address begins with bits 10. It has 16 bit network number (range: 128.0.0.0-to-191.255.255.255), 16 bit host number.
- Class C: Address begins with bits 110. It has 24 bit network number (range: 192.0.0.0-to-223.255.255.255), 8 bit host number.
- Class D: Begins with 1110, multicast addresses (224.0.0.0-to-239.255.255.255)
- Class E: Begins with 11110, unused

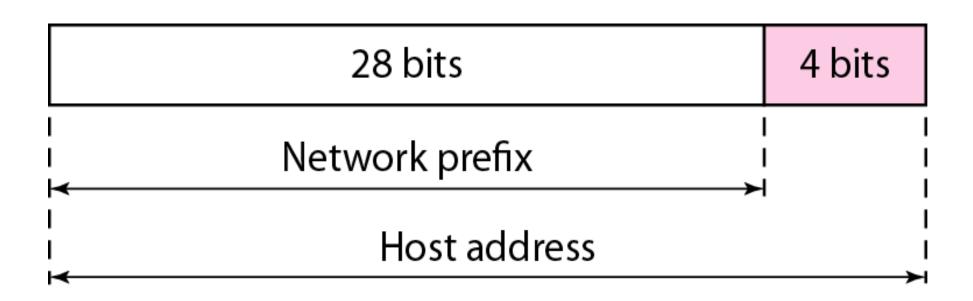
# IPv4 addressing: CIDR

### CIDR: Classless InterDomain Routing

- subnet portion of address of arbitrary length
- address format: a.b.c.d/x, where x is # bits in subnet portion of address



- IPv4 addressing- a block of addresses can be defined as a.b.c.d/x
- a.b.c.d defines one of the addresses and the /x defines the mask.
- Usually a.b.c.d is the first address in the address block.
- The first address in the block can be found by setting the rightmost 32-x bits to 0s.
- The last address in the block can be found by setting the rightmost 32-x bits 1s.
- Number of addresses in the block can be found by using the formula  $2^{32-n}$
- The first address in a block is normally not assigned to any device.
- It is used as the network address that represents the organization to the rest of the world.



 Each address in the block can be considered as a two-level hierarchical structure:

The leftmost n bits (prefix) define the network. The rightmost 32-n bits define the host.

Number of subnet bits is called subnet mask

# Subnetting

• 223.1.17.0/24, ip addresses are  $2^{(32-24)} = 256$ 

• Subnet 1 needs 2^6=64,

223.1.17.0/26

last address: 223.1.17.63

• Subnet 2 needs  $2^6=64$ ,

223.1.17.64/26

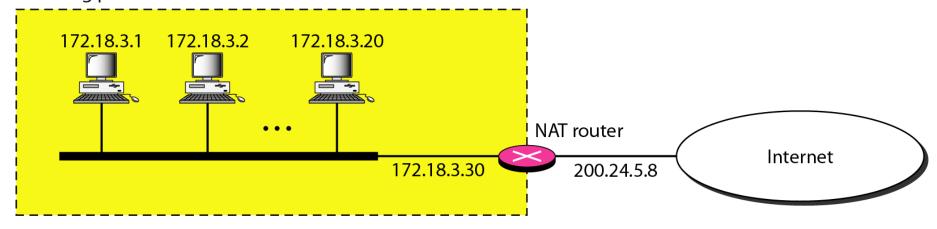
last address: 223.1.17.127

• Subnet 3 needs  $2^7 = 128$ ,

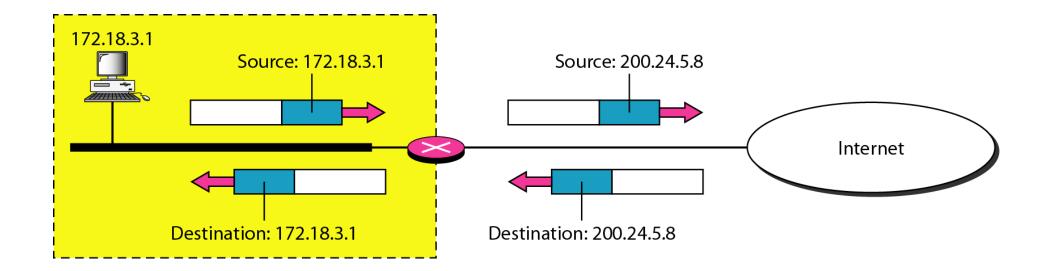
223.1.17.128/25

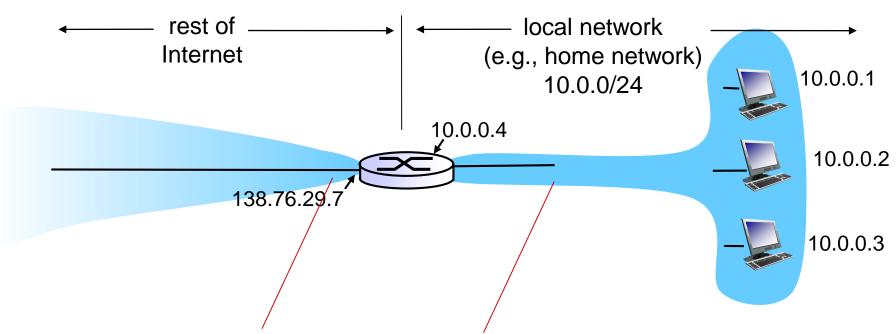
### Figure 19.10 A NAT implementation

#### Site using private addresses



### Figure 19.11 Addresses in a NAT





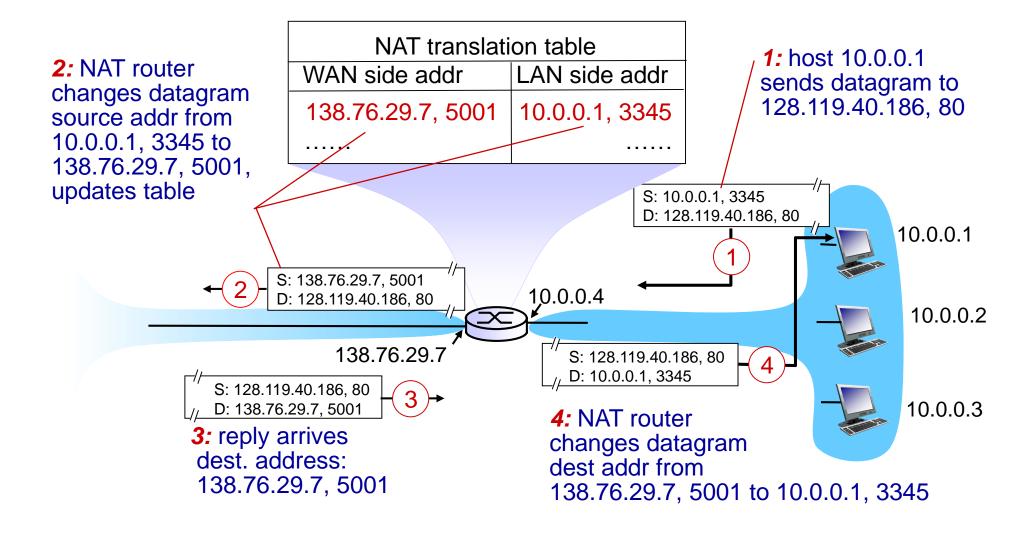
all datagrams leaving local network have same single source NAT IP address: 138.76.29.7, different source port numbers datagrams with source or destination in this network have 10.0.0/24 address for source, destination (as usual)

motivation: local network uses just one IP address as far as outside world is concerned:

- range of addresses not needed from ISP: just one IP address for all devices
- can change addresses of devices in local network without notifying outside world
- can change ISP without changing addresses of devices in local network
- devices inside local net not explicitly addressable, visible by outside world (a security plus)

### implementation: NAT router must:

- outgoing datagrams: replace (source IP address, port #) of every outgoing datagram to (NAT IP address, new port #)
   ... remote clients/servers will respond using (NAT IP address, new port #) as destination addr
- remember (in NAT translation table) every (source IP address, port #) to (NAT IP address, new port #) translation pair
- incoming datagrams: replace (NAT IP address, new port #) in dest fields of every incoming datagram with corresponding (source IP address, port #) stored in NAT table

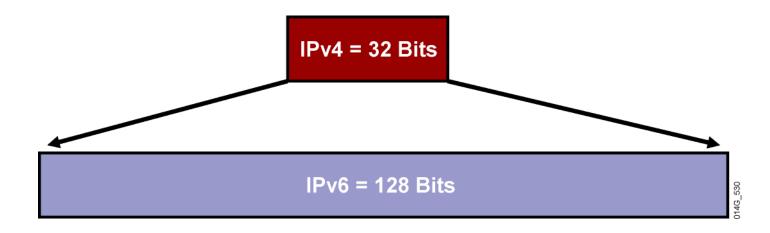


## ICMP: internet control message protocol

*	used by hosts & routers to communicate network-level information	0	0	description echo reply (ping)
	error reporting:	3	0	dest. network unreachable
	unreachable host, network,	3	1	dest host unreachable
		3	2	dest protocol unreachable
	port, protocol	3	3	dest port unreachable
	<ul><li>echo request/reply (used by</li></ul>	3	6	dest network unknown
	ping)	3	7	dest host unknown
*	network-layer "above" IP:	4	0	source quench (congestion
	<ul><li>ICMP msgs carried in IP</li></ul>			control - not used)
	datagrams	8	0	echo request (ping)
		9	0	route advertisement
**	ICMP message: type, code	10	0	router discovery
	plus first 8 bytes of IP	11	0	TTL expired
	datagram causing error	12	0	bad IP header

# **IPv6 Addresses**

- Address depletion may occur with IPv4.
- IPv6 address is 128 bits long.



### **IPv6:** motivation

- initial motivation: 32-bit address space soon to be completely allocated.
- \*additional motivation:
  - header format helps speed processing/forwarding
  - header changes to facilitate QoS

### *IPv6 datagram format:*

- fixed-length 40 byte header
- no fragmentation allowed

# IPv6 datagram format

priority: identify priority among datagrams in flow flow Label: identify datagrams in same "flow." (concept of flow not well defined).

next header: identify upper layer protocol for data

ver	pri	flow label			
K	oayload	llen	next hdr	hop limit	
	source address (128 bits)				
destination address (128 bits)					
data					
← 32 bits — →					

# Other changes from IPv4

- checksum: removed entirely to reduce processing time at each hop
- options: allowed, but outside of header, indicated by "Next Header" field
- \*ICMPv6: new version of ICMP
  - additional message types, e.g. "Packet Too Big"
  - multicast group management functions

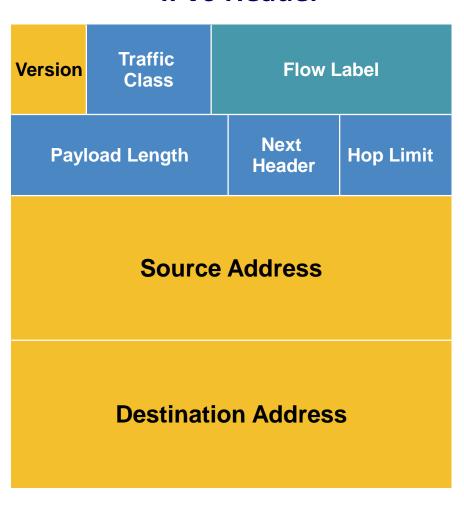
### **IPv6** Header Format

IPv4: 20 Bytes + Options IPv6: 40 Bytes + Extension Header

#### **IPv6** Header

#### **IPv4** Header

Version	IHL	Type of Service	Total Length	
Identification			Flags	Fragment Offset
Time to Live Protocol		Header Checksum		
	Source Address			
Destination Address				
		Options		Padding



# **IPv6 Address Types**

- Unicast
- Address is for a single interface.
- Multicast
- One-to-many
- Enables more efficient use of the network
- Uses a larger address range
- Anycast
  - One-to-nearest (allocated from unicast address space).
  - Multiple devices share the same address.
  - All anycast nodes should provide uniform service.
  - Source devices send packets to anycast address.
  - Routers decide on closest device to reach that destination.
  - Suitable for load balancing and content delivery services.

# IPv6 Address Scope

- Link-local: The scope is the local link (nodes on the same subnet)
- Unique-local: The scope is the organization (private site addressing)
- Global: The scope is global (IPv6 Internet addresses)

• x:x:x:x:x:x:x, where x is a 16-bit hexadecimal field

# IPv6 Address Representation: Link Local

- Hosts on the same link (the same subnet) use these automatically configured addresses to communicate with each other.
- Neighbour discovery provides address resolution.
- Following illustration shows the structure of a link-local address.

1111 1110 10	000 000	Interface ID
(10 bits)	(54 bits)	(64 bits)

# IPv6 Address Representation: Unique Local

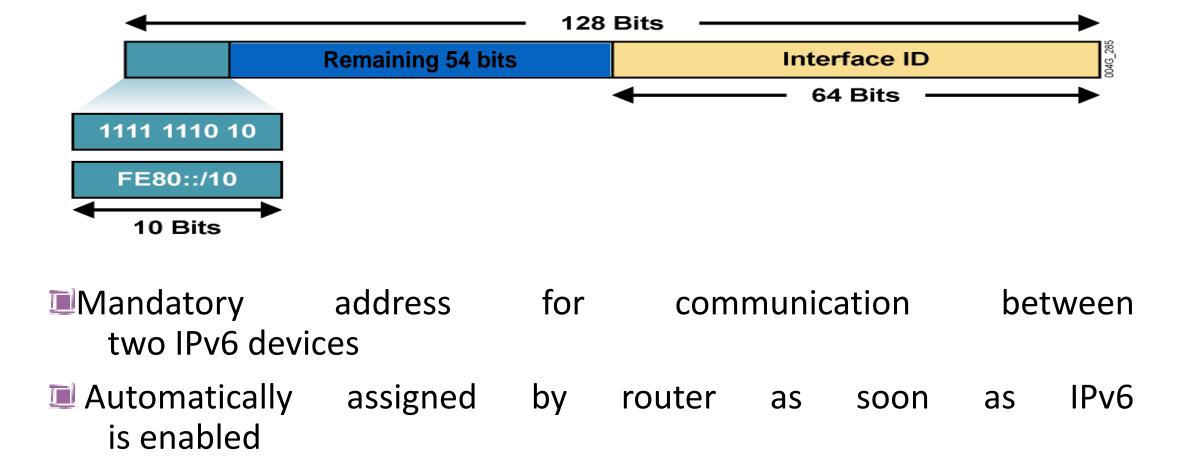
IPv6 unicast unique-local addresses are similar to IPv4 private addresses.

The unique-local address of the scope a internetwork of organization's (You site. an can addresses unique-local global and both use addresses in your network)

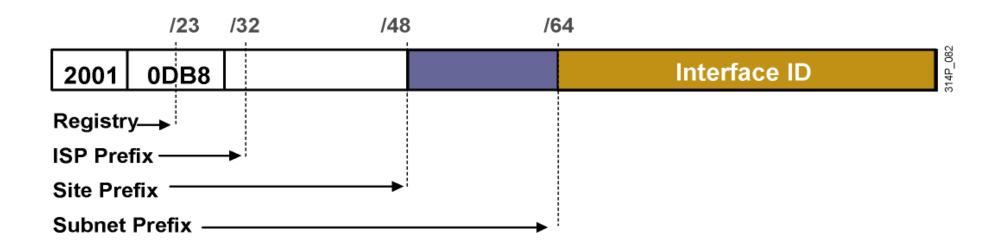
•

The prefix for unique-local addresses is FC00::/8.

# IPv6 Address Representation: Link Local



# IPv6 Address Representation: Global Unicast



Global unicast and anycast addresses are defined by a global routing prefix, a subnet ID, and an interface ID.

### DHCPv6

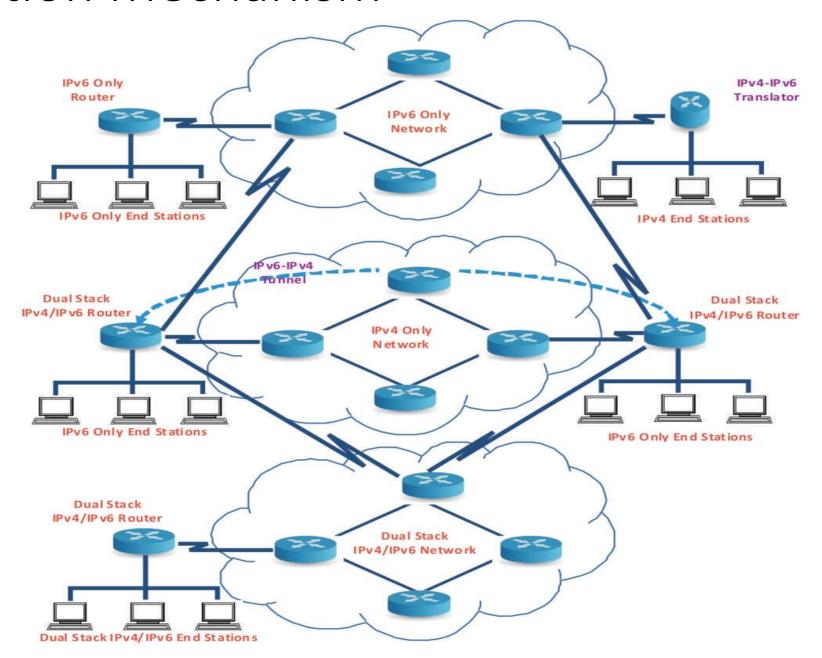
- Stateful Configuration
- Provides not only IP address, also other configuration parameters like DNS

### Transition to IPv6

### **Transition Options:**

- Dual Stack
- IPv6-IPv4 Tunnel
- IPv6-IPv4 Translation

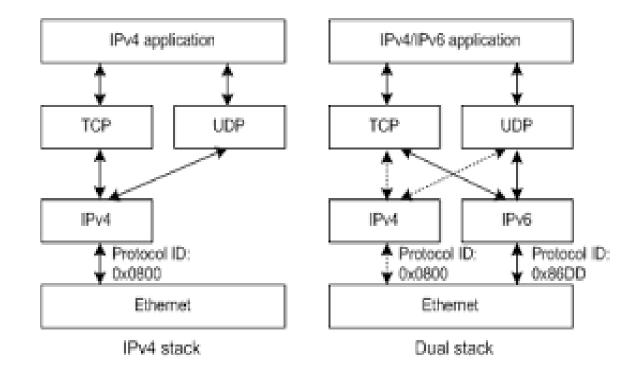
# Transition Mechanism



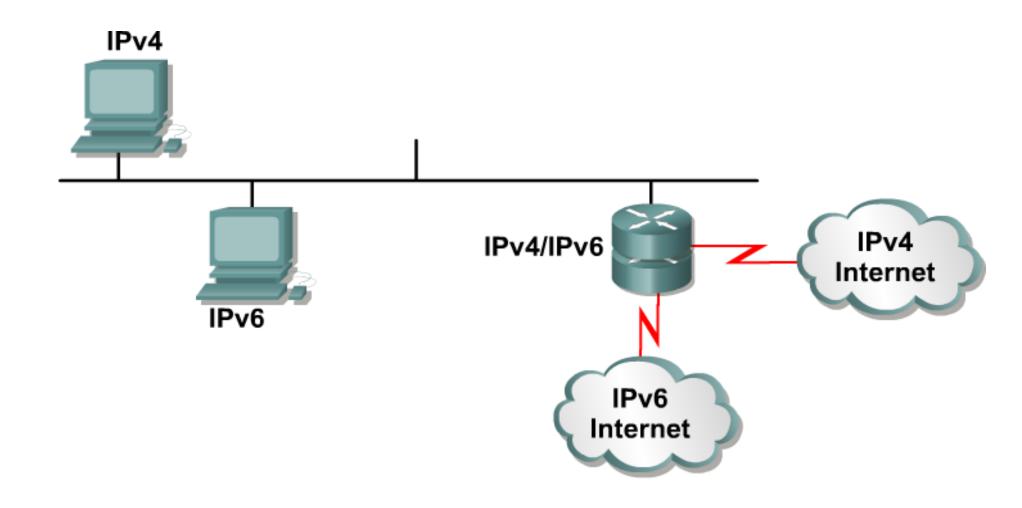
### Dual stack hosts and network

- This allows all the end hosts and intermediate network devices (like routers, switches, modems etc.) to have both IPv4 and IPv6 addresses and protocol stack.
- If both the end stations support IPv6, they can communicate using IPv6; otherwise they will communicate using IPv4.
- This will allow both IPv4 and IPv6 to coexist and slow transition from IPv4 to IPv6 can happen.

## **Dual Stack**

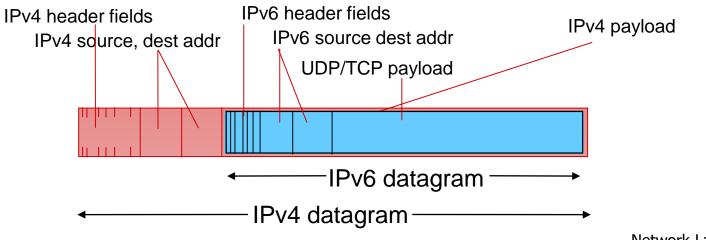


## **Dual Stack Hosts and Network**



### Transition from IPv4 to IPv6

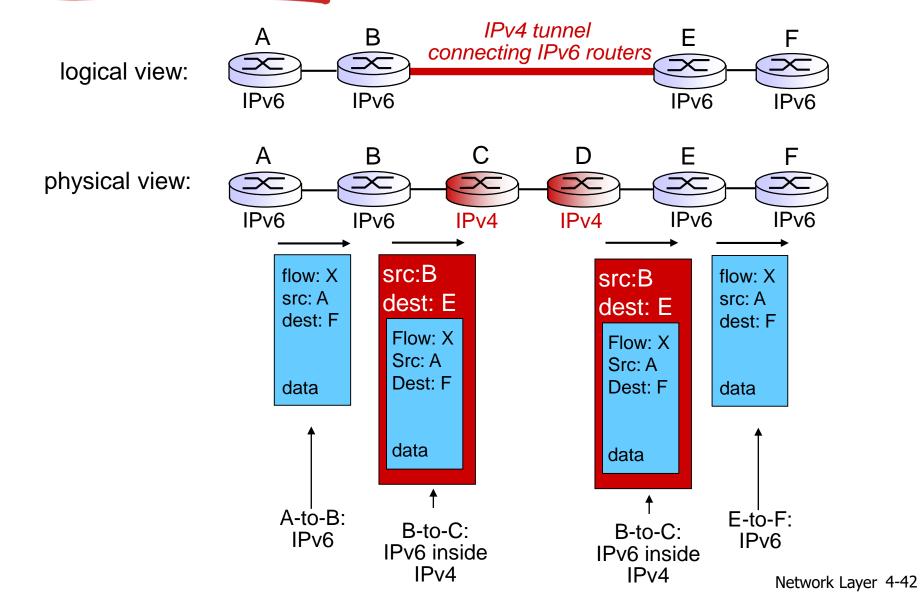
- not all routers can be upgraded simultaneously
  - no "flag days"
  - how will network operate with mixed IPv4 and IPv6 routers?
- tunneling: IPv6 datagram carried as payload in IPv4 datagram among IPv4 routers



# **Tunneling**

IPv4 tunnel Α В connecting IPv6 routers logical view: IPv6 IPv6 IPv6 IPv6 Ε Α В physical view: IPv6 IPv6 IPv6 IPv6 IPv4 IPv4

# **Tunneling**



### **IPv6-IPv4 Translation**

- This allows communication between IPv4 only and IPv6 only end stations
- The job of the translator is to translate IPv6 packets into IPv4 packets by doing address and port translation and vice versa.

