

CS528

Energy Efficient Task Scheduling

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Outline

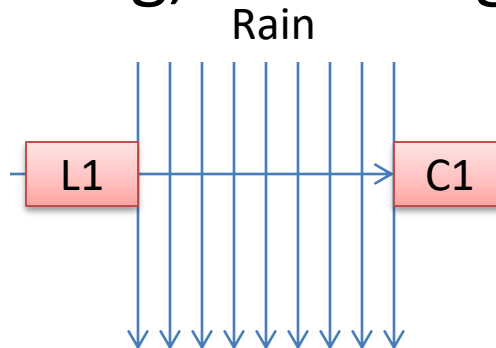
- Power Aware
- Task with Hard Deadlines
- Energy Efficiency
- Energy Efficient Scheduling
- Real Time Tasks

Power Aware Scheduling Vs Energy Aware Scheduling

- Power Budget should not exceed
 - Minimized
 - Monthly Expenses: CAP ==> Solution is EMI
 - Power CAP: If your system have 100 design, at any instance of time you should not run things above 100W
 - Suppose you have 3KW wiring in your home, you have 3 AC with each of 1.5KW rating, At a given time, you can run maximum of 2 AC.
- Total energy budget should not exceed
 - Battery capacity, mah (mobile), AH (UPS)
 - Minimized: EC
 - Power and Time

Speed Matters or Not : I

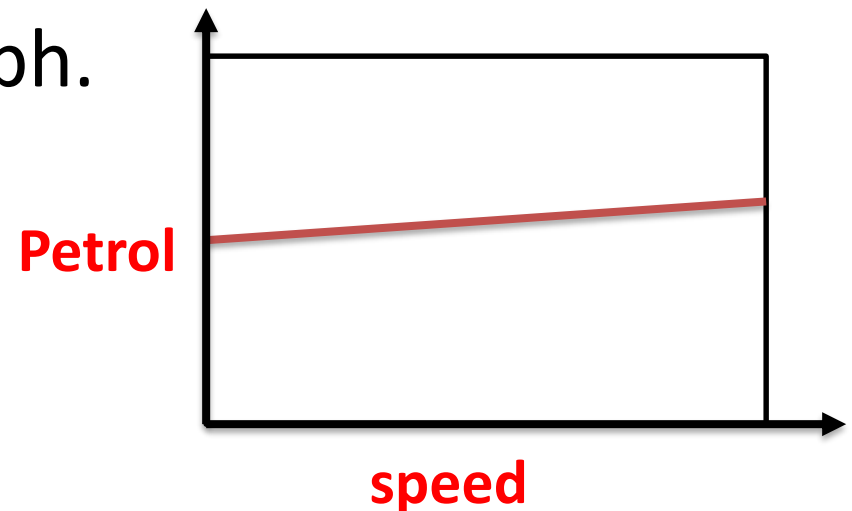
- Assume it is raining, need to go to L1 to C1 of IITG urgently



- Assume rain drops are **falling vertically, uniformly** and you need to walk/run **horizontally**
- Do you get wetter if you run or walk in the rain?**
- Physics answer: Speed does not matter
 - Surface area cover by your body by traveling from L1 to C1 is same, it does not depend on speed

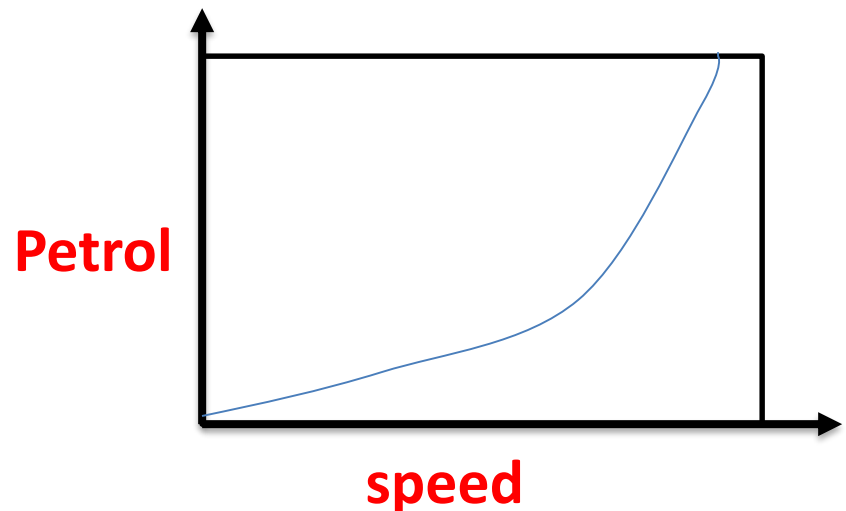
Speed Matters or Not : II

- Assume you have an Royal Enfield Bullet, you need to go from IITG to GS road, 30KM in 1 hours
- Petrol consumption is almost same at any speed. Example it 2ml/minute at 10kmph and 2.1ml/minute at 100kmph.
- How to save petrol ?
 - Sol: Go at higher controllable speed



Speed Matters or Not : II

- Assume you have an Bike, you need to go from IITG to GS road, **30KM in 1 hours**
- Petrol consumption is exponentially/quadratic increasing with speed. Example it 2ml/minute at 10kmph and 20ml/minute at 100kmph.
- How to save petrol ?
 - Sol: Go at slower speed to meet the deadline
 - Above example 30kmph
 - **Critical Speed**



Power and Energy Consumptions

- CPU: dynamic power $P_d = C_{ef} * V_{dd}^2 * f$
 - C_{ef} : switch capacitance, V_{dd} : supply voltage
 - f : processor freq \rightarrow linear related to V_{dd}

$$P \propto f^3$$

- Battery Powered System Reduce Energy usage

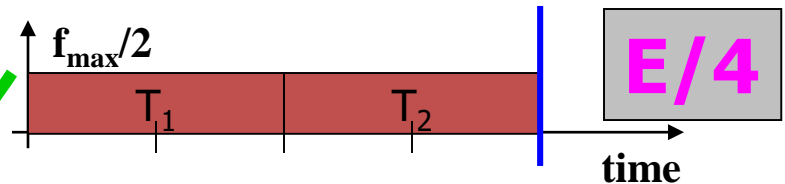
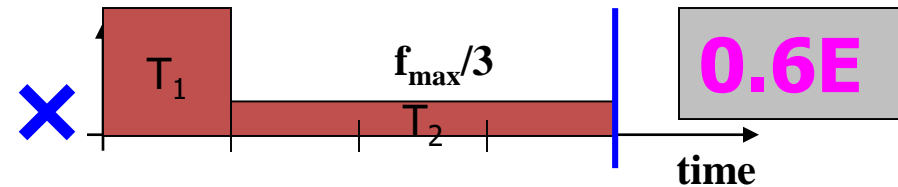
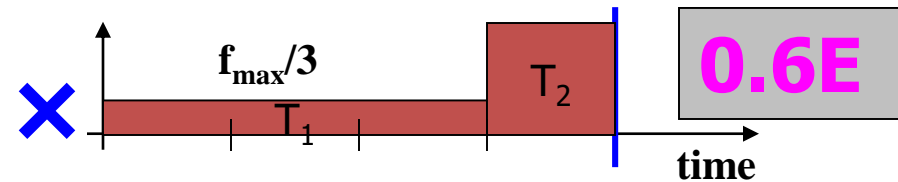
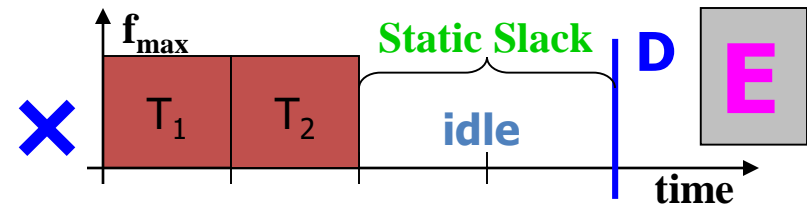
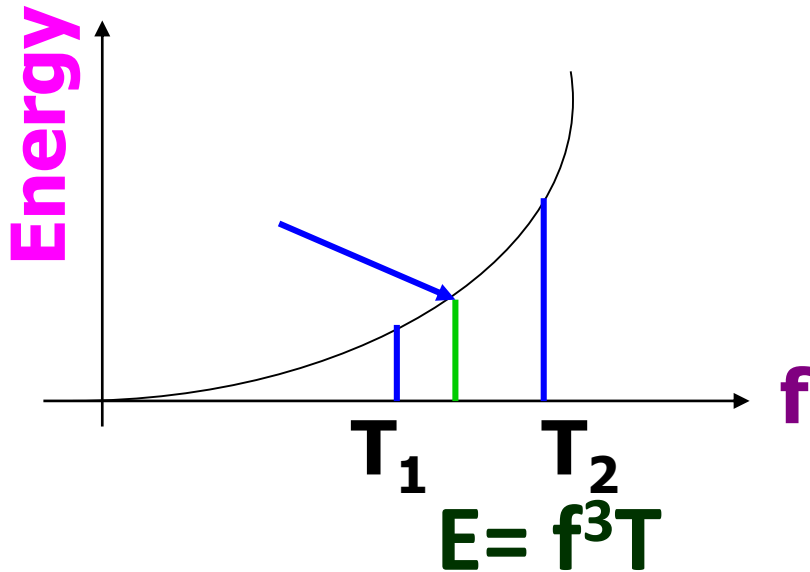
$$E = P \cdot t \propto f^3 t$$

- Execution time t is inverse to f , $t \propto 1/f$

$$\text{So } E \propto f^2$$

Power Aware Scheduling

Static slack: **uniformly** slow down all tasks



I. $f^3 T + f^3 T = 2 \cdot f^3 T = E$

II. $3(f/3)^3 T + f^3 T = 0.57 \cdot E$ ✓

III. $f^3 T + 3(f/3)^3 T = 0.57 \cdot E$

IV. $2(f/2)^3 T + 2(f/2)^3 T = E/4$

Energy Aware Scheduling

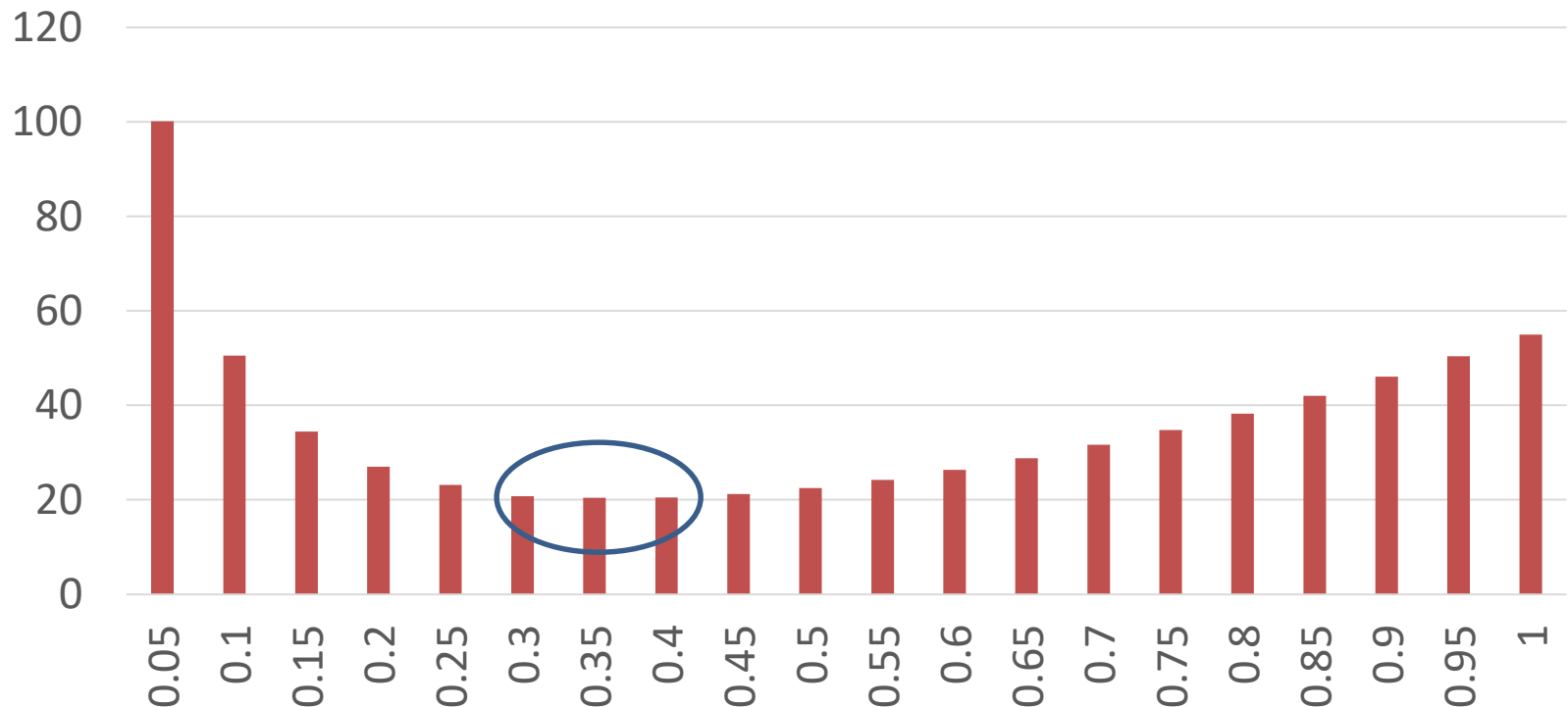
- $E = P * T$, More refined model $P = P_s + k f^3$
- Suppose $f \in [0:1]$.
- $E = (P_s + k f^3) * T / f = T * [P_s / f + k f^2]$
- Min at $dE/df = 0$, $-P_s * 1/(f^2) + 2.k.f = 0$
 $\rightarrow 2kf = P_s/(f^2) \rightarrow f^3 = P_s/2.k$

$$\rightarrow f_c = \sqrt[3]{P_s / 2.k}$$

Full consumption of Splendor

- $P = 5 + 50f^3$
 - Distance to travel 30km in 60 minutes deadline
- $F_c = 0.368399$

Fuel Consumption of Splendor

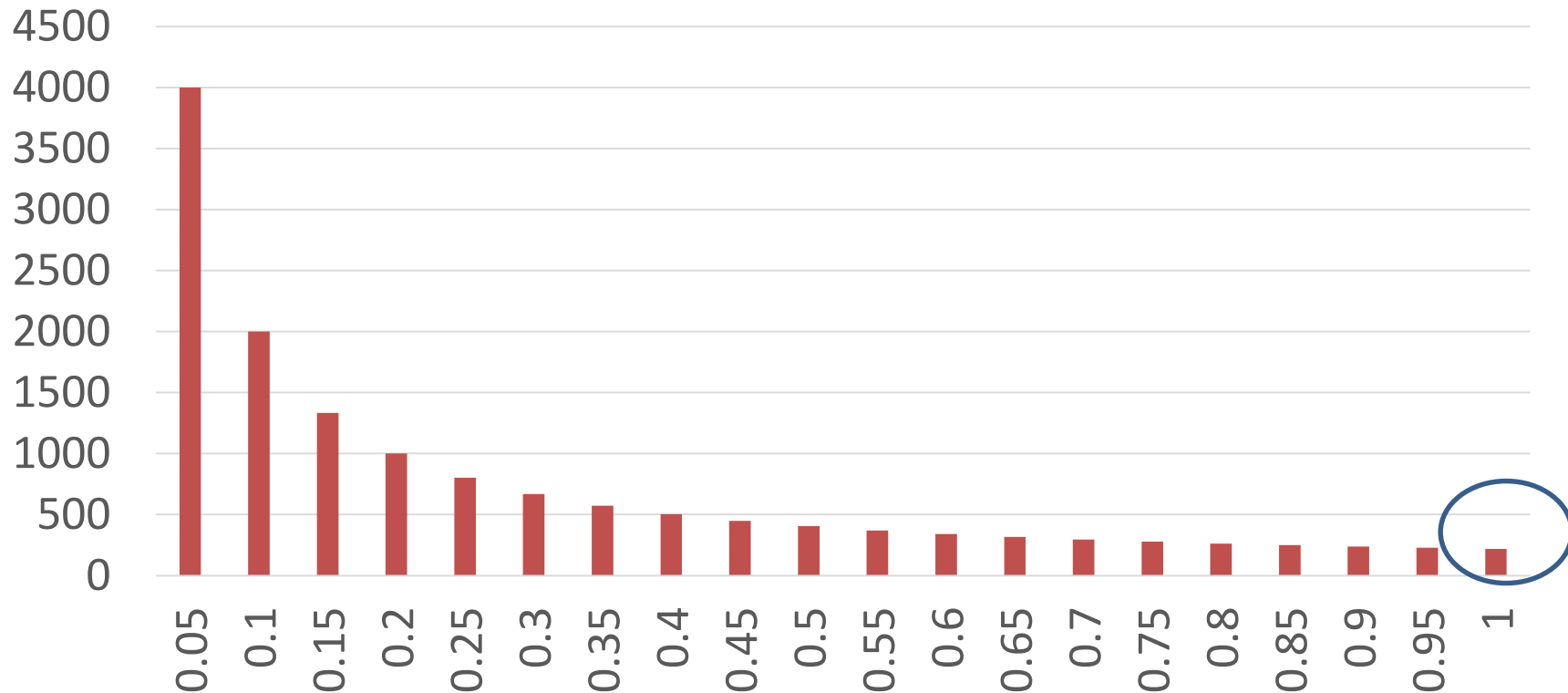


Full consumption of Bullet

- $P = 200 + 20f^3$
- Distance 30km in 60 minutes

– $F_c = 1.709$

Fuel Consumption of Bullet



Problems of Energy Efficiency

- Laptop Problem
 - Given the energy budget, maximize number of Job
 - Given the Budget money maximize your satisfaction
 - Go to Restaurant with Rs 100. Choose Items to fill you stomach with your budget.
 - Given Rs 20 for going from IITG to Airport
 - Go to Jhalukbari using IIT G bus freely, Take another public bus pay Rs 20 to reach Airport.
 - Given Rs 10 : not possible, you need to walk...:)
 - Given Rs 600 how to go : Hire Taxi
 - Given Rs 20000 how to go : Hire BMW/Mercedes along with many other cars for security personals

Problems of Energy Efficiency

- Server Problem
 - Budget is not constraints, minimize budget but do all the work (get all the items)
 - I want to Take all item of Thela/Bora..How much I need to pay? ---Bargaining

Server Problem Example : $P_\infty | p_j, d_j | \Sigma E_j$

- We have infinite processors
- Processor can be run at speed $f=[0:1]$, $PC=\alpha f^3$
- N Tasks with deadlines, Task arrived at time 0, preemption not allowed, p_j at $f=1$
- Execution time task t_j at freq $f = e_j(t_j, f) = p_j/f$;
- Energy consumption task t_j at freq f
 $= E * \text{time} = PC(f) * e_j(t_j, f) = \alpha f^3 p_j / f = \alpha f^2 p_j$
- **We want to execute all the tasks, and minimize the sum of EC of all the tasks**

Server Problem Example : $P_\infty | p_j, d_j | \Sigma E_j$

- We want to execute all the tasks, and minimize the sum of EC of all the tasks
- Solution
 - Select one processor for each of the tasks and total of N processors
 - Run the task at lowest feasible speed to meet the deadline $f_j = p_j / d_j$
- This gives (optimal) minimum ΣE_j
 - Total EC = $\Sigma E_j = \Sigma \alpha f_j^2 p_j$
 - As $(a+b)^2 > a^2 + b^2$: running two task on one processor with higher speed consume higher energy

Laptop Problem Example : $P_\infty, E_b \mid p_j, d_j \mid \Sigma U_j$

- We have infinite processors
- Processor can be run at speed $f=[0:1]$, $PC=\alpha f^3$
- N Tasks with deadlines, Task arrived at time 0, preemption not allowed, p_j at $f=1$
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- **We want to execute maximum number of the tasks before deadline given the energy budget**

Laptop Problem Example : $P_{\infty}, E_b \mid p_j, d_j \mid \sum U_j$

- We want to execute maximum number of the tasks before deadline given the energy budget
- Solution:
 - Sort the tasks based on bare minimum energy requirement $E_j = \alpha f_j^2 p_j$
 - Select the maximum number of task from this set
- Given N item with weight w_1, w_2, \dots, w_N : the weight is critical/min energy required of the task
- Select Maximum number of item given the Budget of Knapsack. **0-1 Knapsack Problem**
- NPC and Pseudo polynomial time algorithm exist using Dynamic Programming.