

Tarea 1

Programacion Declarativa

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1.-

2. a)

$$(f_2 0) \rightarrow_\beta (f_2 3)[(f_2 := \lambda n. \lambda x. \lambda y. yn] = ((\lambda n. \lambda x. \lambda y. yn) 0) \\ \rightarrow_\beta (\lambda x. \lambda y. yn)[n := 0] = (\lambda x. \lambda y. y 0) = 1$$

$$(f_2 3) \rightarrow_\beta (f_2 3)[(f_2 := \lambda n. \lambda x. \lambda y. yn] = ((\lambda n. \lambda x. \lambda y. yn) 3) \\ \rightarrow_\beta (\lambda x. \lambda y. yn)[n := 3] = (\lambda x. \lambda y. y 3) = 4$$

2.b)

$$(g_2 1) \rightarrow_\beta (g_2 1)[(g_2 := \lambda n. n 0 (\lambda x. x))] = ((\lambda n. n 0 (\lambda x. x)) 1) \rightarrow_\beta (n 0 (\lambda x. x))[n := 1] \\ = (n[n := 1] 0[n := 1] (\lambda x. x)[n := 1]) = (1 0 (\lambda x. x)) = ((\lambda x. \lambda y. y 0) 0 (\lambda x. x)) \\ \rightarrow_\beta ((\lambda y. y 0)[x := 0] (\lambda x. x)) = ((\lambda y. y 0) (\lambda x. x)) \rightarrow_\beta (y 0)[y := \lambda x. x] \\ = (y[y := \lambda x. x] 0[y := \lambda x. x]) = ((\lambda x. x) 0) \rightarrow_\beta (x[x := 0]) = 0$$

$$(g_2 4) \rightarrow_\beta (g_2 4)[g_2 := \lambda n. n 0 (\lambda x. x)] = ((\lambda n. n 0 (\lambda x. x)) 4) \rightarrow_\beta (n 0 (\lambda x. x))[n := 4] \\ = (n[n := 4] 0[n := 4] (\lambda x. x)[n := 4]) = (4 0 (\lambda x. x)) = ((\lambda x. \lambda y. y 3) 0 (\lambda x. x)) \\ \rightarrow_\beta ((\lambda y. y 3)[x := 0] (\lambda x. x)) = ((\lambda y. y 3) (\lambda x. x)) \rightarrow_\beta (y 3)[y := \lambda x. x] \\ = (y[y := \lambda x. x] 3[y := \lambda x. x]) = ((\lambda x. x) 3) \rightarrow_\beta (x[x := 3]) = 3$$

2.c)

$$(h_2 0) \rightarrow_\beta (h_2 0)[(h_2 := \lambda n. n \underline{true} (\lambda x. \underline{false})] = ((\lambda n. n \underline{true} (\lambda x. \underline{false})) 0) \\ \rightarrow_\beta (n \underline{true} (\lambda x. \underline{false}))[n := 0] = (n[n := 0] \underline{true}[n := 0] (\lambda x. \underline{false}))[n := 0] \\ = (0 \underline{true} (\lambda x. \underline{false})) = ((\lambda x. \lambda y. x) \underline{true} (\lambda x. \underline{false})) \rightarrow_\beta ((\lambda y. x) [x := \underline{true}] (\lambda x. \underline{false}))$$

$$= ((\lambda y.\underline{true}) (\lambda x.\underline{false})) = \underline{true}$$

$$\begin{aligned} (h_2 \ 5) &\rightarrow_\beta (h_2 \ 5)[(h_2 := \lambda n.n \ \underline{true} \ (\lambda x.\underline{false}))] = ((\lambda n.n \ \underline{true} \ (\lambda x.\underline{false})) \ 5) \\ &\rightarrow_\beta (n \ \underline{true} \ (\lambda x.\underline{false}))[n := 5] = (n[n := 5] \ \underline{true}[n := 5] \ (\lambda x.\underline{false}))[n := 5] \\ &= (5 \ \underline{true} \ (\lambda x.\underline{false})) = ((\lambda x.\lambda y.y \ 4) \ \underline{true} \ (\lambda x.\underline{false})) \rightarrow_\beta ((\lambda y.y \ 4) [x := \underline{true}] \ (\lambda x.\underline{false})) \\ &= ((\lambda y.y \ 4) (\lambda x.\underline{false})) \rightarrow_\beta ((y \ 4)[y := (\lambda x.\underline{false})]) = (\lambda x.\underline{false} \ 4) = \underline{false} \end{aligned}$$

2.d)

f_2 es una función que suma 1 al numero pasado como argumento (el numero tiene que ser un natural de Scott)

g_2 es una función que resta 1 al numero pasado como argumento (el numero tiene que ser un natural de Scott)

h_2 es una función que nos dice si el argumento es 0 el argumento es un numero natural de Scott

2.e)

tenemos que :

$$\underline{true} = \lambda x.\lambda y.x$$

$$\underline{false} = \lambda x.\lambda y.y$$

$$\underline{if} = \lambda p.\lambda a.\lambda b \ p \ a \ b$$

$$\underline{h_2} = \lambda n.n \ \underline{true} \ (\lambda x.\underline{false})$$

$$\underline{f_2} = \lambda n.\lambda x.\lambda y.yn$$

$$\underline{g_2} = \lambda n.n \ 0 \ (\lambda x.x)$$

Sea F un combinador de punto fijo:

Definimos $\underline{sumaScott} \Rightarrow Fg$ donde

$$g \Rightarrow \lambda f.\lambda x.\lambda y.\underline{if} \ (\underline{h_2} \ y) \ x \ (f \ (\underline{f_2} \ x) \ (\underline{g_2} \ y))$$