

//Problem 1

//1.1)

```
int minimumCost(int *size, float *cost, int k, int n) {

    float best[n+1] = infinity //set all values to infinity
    best[0] = 0.0
    for (int i = 1; i < n+1; ++i)
    {
        for (int j = 0; j < k; ++j)
        {
            if ((i - size[j]) < 0) {
                continue;
            }

            best[i] = min(best[i],best[i-size[j]] + cost[j])
        }
    }
    return best[n]
}
```

//1.2)

//this function assumes there is a way to make n pierogies

```
int *sizes(int *OPT, int *size, float *cost, int k, int n) {
    int *usedList = malloc(sizeof(int) * (n+1)) //won't be this large but just in case
    int used = 0

    int i = n
    while (i > 0) {
        for (int j = 0; j < k; ++j)
        {
            if ((i - size[j]) >= 0) {
                if (OPT[i] >= (OPT[i-size[j]] + cost[j])) {
                    usedList[used] = size[j]
                    used++
                    i = i - size[j]
                    break; //breaks out of the for loop. no need to finish all k
checks.
                }
            }
        }
    }
    usedList = realloc(sizeof(int) * (used+1))
    return usedList
}
```

/*

For part 1: function minimumCost

$O(n*k)$ run time:

This can be seen clearly with the for loop inside another for loop.

i goes from 1 to n and j goes from 0 to k-1.

$O(n)$ space:

array best takes $n+1$ space.

For part 2: function sizes

$O(nk)$ run time. $O(n)$ space.

i in its worst case goes from n to 0.

But in any average case it will skip many of the numbers from n to 0 because of $i = i - \text{size}[j]$.

j always goes from 0 to k. This leads to $O(n*k)$ run time.

usedList takes worst case size of $n+1$ but in most cases it's much less.

*/

//Problem 2

```
int maxnetvalue(int *v, int n, int m) {
    int n //length of rod
    int m //cost per cut
    int v[n+1] //value per length v[1] to v[n] are inputted values
    int best[n+1] //will calculate best value for length 0 to length n

    best[0] = 0 //base case: best value for length 0 is 0
    for (int i = 1; i < n+1; ++i)
    {
        best[i] = v[i]
    }

    for (int length = 3; length < n+1; ++length)
    {
        int numPairs = (int) ((length-1)/2)+1
        for (int j = 1; j < numPairs; ++j)
        {
            best[length] = max(best[length], v[length-1-j]+v[j]-m)
            //should have figured out lower cases with optimal cuts so won't
            need to evaluate more than one cut
        }
    }

    return best[n]
}
/*
O(n^2) run time. This can be seen clearly with the for loop inside another for loop.
length can max be size n and j can max reach size (n-1)/2. n * (n-1)/2 is O(n^2) time.
initialization takes O(n) time
*/
```

//Problem 3

```
float averageWeight(int **edge, int n) { //edge is an adjacency matrix
    int cost[end+1] //really only needs to go from start to end. But it's easier to index this
way.
    int numPaths[end+1]

    //base cases:
    numPaths[n-1] = 1
    cost[n-1] = 0

    //initialize everything else
    for (int i = 0; i < n - 1; ++i)
    {
        numPaths[i] = 0
        cost[i] = 0
    }

    for (int i = n - 1; i >= 0; --i)
    {
        for (int j = i + 1; j < n + 1; ++j)
        {
            if (edge[i][j] != 0) {
                numPaths[i] += numPaths[j]
                cost[i] += (cost[j] + 1*numPaths[j]) //all edges are unweighted and
equal to 1
            }
        }
    }

    if (numPaths[0] == 0) {
        return NaN
    }
    else {
        return cost[0]/numPaths[0]
    }
}

/*
O(n^2) run time. This can be seen clearly with the for loop inside another for loop.
i is size n going from n - 1 to 0 and j can max reach size n when i = 0.
initialization takes O(n) time
*/
```