

Setup

- Generating Random Text String:
 - I used Java's Random class to generate a random string of 10,000 characters that included alphabets, numbers, and special characters.
 - o Code: The function takes in the length of the random string desired, generates and returns it.

• Generating Random Pattern Strings:

- For each m, I generated 1000 random patterns using the same function above. Then calculated the average running time for finding each pattern in the text.
 - Eg: m = 2 -> ig, 5j, !p. AB, Kp, (1000 random patterns)
 - Eg: m = 3 -> oie, 19f, ;rt, qec, (1000 random patterns)
- Since these pattern strings were generated at random, some were present in the Text string, and some weren't.
- Since the code tested the execution times for 1000 random patterns, the average elapsed time
 was a good indicator of the real-world execution time of each algorithm.
- Reading speeches.txt file for English Words Matching
 - I downloaded and stored the speeches.txt file in the same directory as the code.
 - o The code used Java's Files.readAllBytes() function

```
byte[] bytes = Files.readAllBytes(Paths.get("project2/speeches.txt"));
String text = new String(bytes);
```

- Reading words from words1.txt file for English Words Matching
 - I used the same words1.txt file from Assignment 1.
 - o The code read the words using BufferedReader and stored the words in a HashMap of length to List of words. Screenshot from IntelliJ:

```
Result:

V OO result = {HashMap@837} size = 11

V = {Integer@852} 2 -> {ArrayList@853} size = 1271

| key = {Integer@852} 2
| value = {ArrayList@853} "[2D, 3D, 3M, 4H, -a, A., a', a-, a., A1, a1, A4, A5, AA, aa, AB, Ab, ab, AC, Ac, ac, AD, Ad, ad, AE... View
| v = {Integer@854} 3 -> {ArrayList@855} size = 6221
| key = {Integer@854} 3
| value = {ArrayList@855} "[1st, 2nd, 3-D, 3-d, 3rd, 4-D, 4GL, 4th, 5-T, 5th, 6th, 7th, 8th, 9th, A-1, AAA, aaa, AAE, AAF, AA... View
| lnteger@856} 4 -> {ArrayList@857} size = 13208
```

- From each M above, the code randomly samples 1000 words from the list and calculates the average over multiple runs.
- o Code:

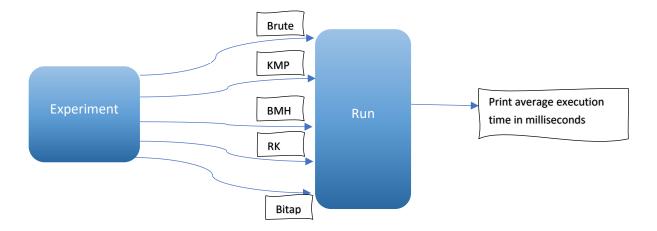
```
Random random = new Random();
for (int 1 : lengths) {
   int numberOfWords = wordMap.get(l).size();
   for (int i = 0; i < 1000; i++) {
      m[i] = wordMap.get(l).get(random.nextInt(numberOfWords)); // choose a
   random word from list
   }
   experiment(m[0].length(), m, text); // experiment 1000 times for each m
}</pre>
```

- Calculating the Execution Time
 - o lused Java's System.nanoTime()
 - Code: The function takes in the algorithm to run. Executes the algorithm for the given 1000 pattern strings. And prints the average execution time.

```
private static void run(String[] M, String T, String algo, Matcher matcher) {
    double totalElapsed = 0;
    for (String P : M) {
        long start = System.nanoTime();
        int res = matcher.match(T, P);
        long finish = System.nanoTime();
        totalElapsed += finish - start;
    }
    totalElapsed /= 1000.0;
    System.out.printf(algo +": %.5f microseconds", totalElapsed / 1000.0);
}
```

 The Above run method is called for all the five algorithms by passing the desired algorithm object. All the algorithm class extend the Matcher class and use inheritance

```
private static void experiment(int mLength, String[] M, String T) {
    run(M, T, "Brute", brute);
    run(M, T, "KMP ", kmp);
    run(M, T, "BMH ", bmh);
    run(M, T, "RK ", rk);
    run(M, T, "BITAP", bitap);
}
```



Pattern	Time in microseconds					
m	Brute Force	KMP	вмн	RK	BITAP	
2	32.81333	19.77967	33.70487	45.4758	27.41065	
3	25.31113	11.50986	14.28385	40.6214	21.51133	
4	25.52351	11.7321	10.53296	41.37416	21.68629	
5	25.74828	11.56085	8.83097	40.95325	21.70076	
6	26.31965	11.31291	7.48713	40.91029	22.04134	
7	26.34072	11.60992	6.22689	40.90109	21.60537	
8	25.52076	11.53245	7.10483	42.05945	21.9666	
9	25.91869	11.27961	5.43424	41.07827	21.83601	
10	25.7609	11.71329	4.50006	40.75521	21.6186	

Table 1. Execution times of random strings

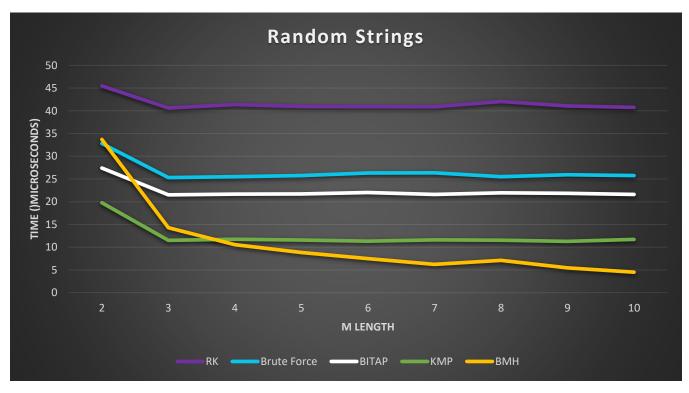


Fig 1. Plot of execution times of random strings

Pattern	Time in microseconds					
m	Brute Force	KMP	ВМН	RK	BITAP	
2	79.57688	36.89431	78.64268	129.74005	97.50817	
3	98.17098	62.53191	58.61704	162.82624	96.33735	
4	110.33659	58.61114	50.0703	170.57134	101.65993	
5	115.41107	55.16546	44.38198	178.45565	106.15033	
6	120.03385	65.71231	39.0368	179.27862	107.12553	
7	121.74034	83.18445	34.50343	180.43071	107.2616	
8	123.32408	56.76527	31.17611	179.86031	107.34801	
9	123.86736	58.00938	28.71236	180.05006	107.41769	
10	125.4105	60.30199	26.78517	180.39783	107.59897	

Table 2. Execution times of English words

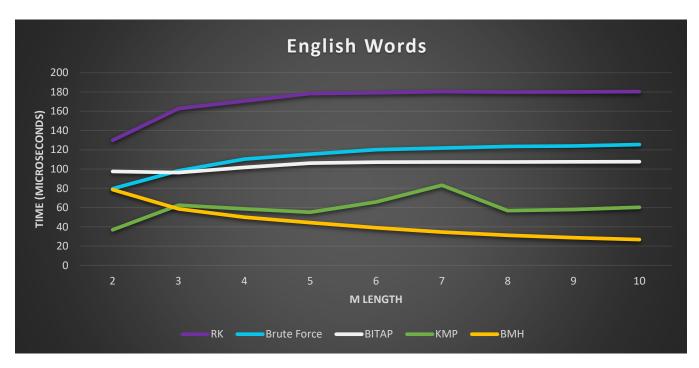


Fig 2. Plot of execution times of English words

Conclusions and observations

- From the above plots, it is clear that the best algorithm for both random strings and English words is **Boyer-Moore-Harspool**. The **trendline for BMH** shows that it steadily decreases as the length of m increases. This is highly desirable since in real-world, extremely small patterns are unlikely.
- Only for extremely small pattern length 1 or 2, it can be seen that BMH falls behind other algorithms. But for other lengths it is better than almost every other algorithm. This is due to the fact that the Last Occurrence Table would mostly be empty for the Alphabet size, and it essentially becomes sliding P across T. Therefore, it tries to match almost every single character of T with P.
- Furthermore, my empirical results show that Rabin-Karp performs the worst. This is surprising but understandable because even though the average and best-case complexity of Rabin-Karp is O(m+n), its worst-case is O(mn) and is inferior to other algorithms. Furthermore, it is noteworthy to point out that it performs worse than brute force this could be due to the overhead of Cyclic Polynomial Rolling Hash and Bit manipulation. If we have multiple patterns to match, a modified Rabin-Karp algorithm might perform better.
- Both **KMP** and **Bitap** perform better than **brute force**. However, if our pattern size tends to be larger than the **word length** of the machine, KMP would significantly outperform Bitap. Bitap also prefers significantly smaller Alphabet size.
- Interestingly, even though theoretical asymptotic efficiencies hint that KMP would be the best performer, in real-world scenarios **constant factors add-up very quickly**. Linear worst-case algorithms such as KMP have a significant overhead of pre-processing the entire Pattern String. For this reason, **KMP is also not space optimal for large P**.
- Unlike KMP, BMH takes advantage of the fact that in real-word scenarios Alphabet size is considerably limited — hence preprocessing is faster for longer Pattern Strings. It is also not limited by the word length of the machine like Bitap. Making it the best algorithm for our scenario.

Algorithm	Preprocessing required?	Performs best when	Performs worst when
Brute	No	Pattern present early on	Pattern repeatedly
			present with character
			mismatch towards the end
Rabin-Karp	Yes	Multi pattern search	Hash collision/overhead
			and limited size of integer
			data type size
KMP	Yes	Pattern significantly	Large Pattern and large
		smaller than alphabet	alphabet (more
		size. Example: Large files	mismatches)
Bitap	Yes	Significantly small	Alphabet size increases
		alphabet size	
Boyer-Moore-Harspool	Yes	Large pattern length and	Low bad character skip
		small alphabet size	

Screenshots from openlab.ics.uci.edu

Random Strings:

```
$ java project2/Runner
Executing Random String matches...
Brute: 33.09176 microseconds
KMP : 20.14450 microseconds
BMH : 33.61610 microseconds
     : 42.85427 microseconds
BITAP: 27.65414 microseconds
Brute: 25.21609 microseconds
KMP : 11.22853 microseconds
BMH : 14.05341 microseconds
      : 40.47439 microseconds
BITAP: 21.42560 microseconds
Brute: 25.68704 microseconds
KMP : 11.68417 microseconds
BMH : 10.54101 microseconds
      : 40.80869 microseconds
BITAP: 21.68672 microseconds
Brute: 25.68794 microseconds
KMP : 11.66298 microseconds
BMH : 8.82303 microseconds
      : 40.74316 microseconds
BITAP: 21.69177 microseconds
Brute: 25.68452 microseconds
KMP : 11.40239 microseconds
BMH : 7.48916 microseconds
      : 40.70956 microseconds
BITAP: 21.72070 microseconds
Brute: 25.68279 microseconds
KMP : 11.71891 microseconds
BMH : 6.21370 microseconds
      : 40.69392 microseconds
BITAP: 21.59305 microseconds
Brute: 25.51855 microseconds
KMP : 11.63090 microseconds
BMH : 5.49012 microseconds
      : 40.85653 microseconds
BITAP: 21.57927 microseconds
Brute: 25.51051 microseconds
KMP : 11.34425 microseconds
BMH : 5.39208 microseconds
     : 40.93893 microseconds
BITAP: 21.58491 microseconds
Brute: 25.49380 microseconds
KMP : 11.65576 microseconds
BMH : 4.48656 microseconds
      : 40.85457 microseconds
BITAP: 21.59257 microseconds
```

• English Words:

```
Executing English word matches...
M = 2
Brute: 81.13885 microseconds
KMP : 37.00685 microseconds
BMH : 80.63475 microseconds
RK : 135.39251 microseconds
BITAP: 96.25218 microseconds
M = 3
Brute: 98.43857 microseconds
KMP : 64.85531 microseconds
BMH : 56.20978 microseconds
     : 160.97909 microseconds
BITAP: 95.31727 microseconds
M = 4
Brute: 113.18915 microseconds
KMP : 45.33125 microseconds
BMH : 49.27327 microseconds
RK : 168.77698 microseconds
BITAP: 100.52441 microseconds
Brute: 116.94712 microseconds
KMP : 48.84299 microseconds
BMH : 44.74250 microseconds
RK : 178.29368 microseconds
BITAP: 107.10072 microseconds
Brute: 119.84806 microseconds
HIP: 19.04000 MILCOSECONDS

KMP : 52.10202 microseconds

BMH : 39.06986 microseconds

RK : 178.79147 microseconds

BITAP: 107.01682 microseconds
Brute: 121.78047 microseconds
KMP : 77.95015 microseconds
BMH : 35.05941 microseconds
RK : 180.23182 microseconds
BITAP: 107.29925 microseconds
M = 8
Brute: 123.66076 microseconds
KMP : 62.67563 microseconds
BMH : 31.58803 microseconds
RK : 181.33488 microseconds
BITAP: 107.89859 microseconds
M = 9
Brute: 123.88495 microseconds
KMP : 58.45241 microseconds
BMH : 28.70399 microseconds
RK : 180.34961 microseconds
BITAP: 107.60400 microseconds
M = 10
Brute: 125.67780 microseconds
KMP : 60.94490 microseconds
BMH : 26.50869 microseconds
      : 180.04053 microseconds
BITAP: 107.54691 microseconds
```