# Data Structures Circular Queue

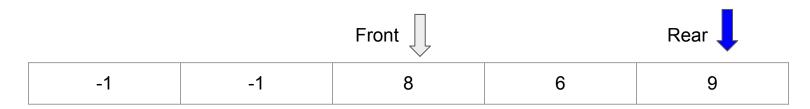
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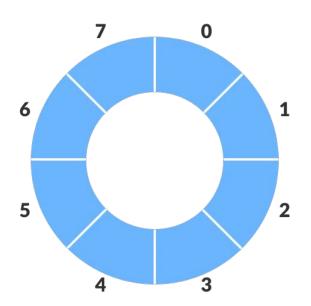
# Array-based: Front-Read approach

- We will have 2 indices: front and rear representing start to end in array
  - When we enqueue element we add it in rear
  - $\circ$  When we dequeue element we shift front to the right  $\Rightarrow$  O(1)
- Enqueue 3: ERROR Queue is full!
- Wait, but there are slots empty in the begin!
  - This is a critical drawback in this approach
  - How to solve? Think for 15 minutes!



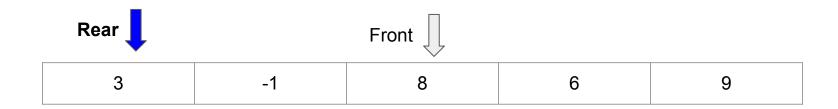
### Circular Queue

- There is a simple way to solve the previous space issue
- Simply, think in the array as a circle
  - o On right, an array of 8 elements, as a circle
- That is, after the last element, there is actually another element, which is position 0
- Now, the queue is full IFF all elements are in use



# Array-based Circular queue

- We will have 2 indices: front and rear representing start to end in array
  - When we enqueue element we add it in rear
  - When we dequeue element we shift font to the right  $\Rightarrow$  O(1)
- Enqueue 3
  - Now move from the last index to index 0 and put the new element
  - Observe: Rear now is BEFORE front



### Initial values for rear & front

- There are several approaches for that
  - In every approach, we have to be consistent in the whole implementation
  - Careful conditions for IsEmpty and IsFull
- Possible initializations

```
rear = front = -1 [initially equal]
rear = front = 0 [initially equal]
rear = -1 and front = 0 [initially !equal]
rear = size - 1 and front = 0 [initially !equal]
```

### int added\_elements = 0;

To avoid tricky conditions and simplify coding, maintain counter for number of elements!

### Circular Queue: Data Structure

- We initially use front = rear = 0
  - o To add a new element, **add in rear** and **move** rear
  - To dequeue element, get front and move front

```
class Queue {
   int size { };
   int front { 0 };
   int rear { 0 };
   int added_elements { };
   int *array { };
}
```

### Circular Queue: Move index

- To move an index step forward consider
  - If this is last element in the array, next position = 0
  - We can do that with if condition (efficient)
  - Or with simple mod

-1

Assume size = 5. Let's try positions from 0 to 5

-1

```
0 % 5 = 0
```

$$0 5 \% 5 = 0$$

```
int next(int pos) {
    //return (pos + 1) % size;

++pos;
    if (pos == size)
        pos = 0;
    return pos;
}
```

- Initially an empty queue. Both rear = front = 0
- Observe: Empty queue with rear == front



• Enqueue (1)  $\Rightarrow$  Add in rear position and move it

Front	Rear		
1			

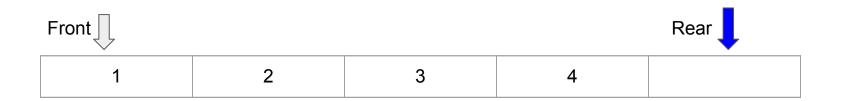
• Enqueue (1), Enqueue (2)



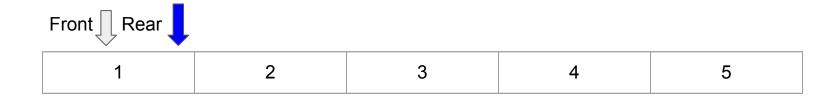
• Enqueue (1), Enqueue (2), Enqueue (3)



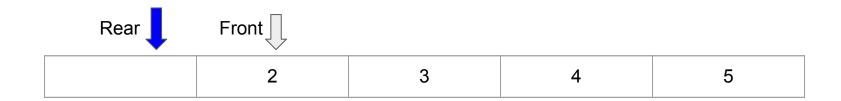
- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4)
- Observe: rear at last array position
  - One more enqueue and it moves the index to 0



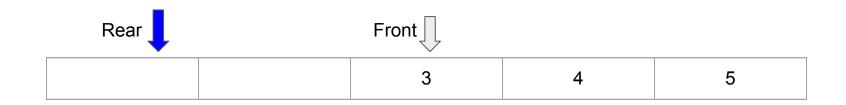
- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4), Enqueue (5)
- Observe: full queue but also rear == front
  - How can we know array is empty or full this way!
  - o We can't!
  - Use the added elements variable
    - 0 = empty
    - 5 = full



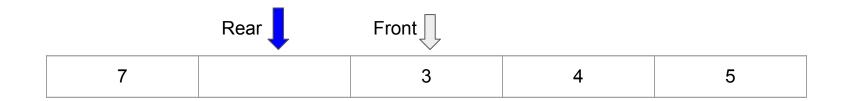
- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4), Enqueue (5)
- Dequeue ⇒ 1
- Observe: Front after Rear



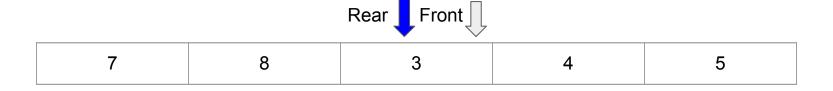
- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4), Enqueue (5)
- Dequeue ⇒ 1
- Dequeue  $\Rightarrow$  2



- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4), Enqueue (5)
- Dequeue ⇒ 1
- Dequeue  $\Rightarrow$  2
- Enqueue (7)



- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4), Enqueue (5)
- Dequeue ⇒ 1
- Dequeue ⇒ 2
- Enqueue (7), Enqueue (8)
- Again full but rear = front = index 2



- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4), Enqueue (5)
- Dequeue ⇒ 1
- Dequeue  $\Rightarrow$  2
- Enqueue (7), Enqueue (8)
- Dequeue  $\Rightarrow$  3



- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4), Enqueue (5)
- Dequeue ⇒ 1
- Dequeue ⇒ 2
- Enqueue (7), Enqueue (8)
- Dequeue  $\Rightarrow$  3
- Dequeue ⇒ 4



- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4), Enqueue (5)
- Dequeue ⇒ 1
- Dequeue ⇒ 2
- Enqueue (7), Enqueue (8)
- Dequeue  $\Rightarrow$  3
- Dequeue ⇒ 4
- Dequeue, Dequeue, Dequeue ⇒ 5, 7, 8
  - Observe: empty with front = rear = 2



# IsEmpty? IsFull?

Trivially handled using added\_elements

```
int isEmpty() {
    return added_elements == 0;
}

bool isFull() {
    return added_elements == size;
}
```

# Enqueue and Dequeue

- Enqueue: Add in rear and move
- Dequeue: Get from front and move
- Direct!

```
void enqueue(int value) {
    assert(!isFull());
    array[rear] = value;
    rear = next(rear);
    added elements++;
int dequeue() {
    assert(!isEmpty());
    int value = array[front];
    front = next(front);
    --added elements;
    return value;
```

# Display Queue

Simply start from the front and count based on added\_elements

```
void display() {
    cout << "Front " << front << " - rear " << rear << "\t";
    if (isFull())
        cout << "full":
    else if (isEmpty()) {
        cout << "empty\n\n";</pre>
        return;
    cout << "\n";
    for (int cur = front, step = 0; step < added elements;</pre>
            ++step, cur = next(cur))
        cout << array[cur] << " ";
    cout << "\n\n";
```

"Acquire knowledge and impart it to the people."

"Seek knowledge from the Cradle to the Grave."