

# *Data Structures*

## Appending operation

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# Get data

- 2 nice utilities might be to get the front or the back element
- But what if we want to:
  - Add in the back?
    - We call it append operation
    - Or push\_back operation
  - Similarly insert in the front?
    - We call it push\_front
- Take 10 minutes:
  - Create function void **push\_back**(int value)
  - It adds an element to the current array

```
int get_front() {  
    return arr[0];  
}  
  
int get_back() {  
    return arr[size-1];  
}
```

# push\_back

- To add an element to an array, we need each time
  - Create new array
  - Move old data to it
  - Add new value
  - Use the new data and remove old one

```
void push_back(int value) {  
    // Add element to the end of the vector  
    // 1) create a new array of bigger size  
    int *arr2 = new int[size + 1];  
    // 2) copy old data  
    for (int i = 0; i < size; ++i)  
        arr2[i] = arr[i];  
    // 3) add the new element and increase size  
    arr2[size++] = value;  
    // 4) change the pointers  
    swap(arr, arr2);  
    // 5) remove the useless data  
    delete[] arr2;  
}
```

# Usage

- One clear advantage now compare to primitive arrays, our data structure can **grows normally in size!**
- From **efficiency** perspective, what is wrong with our code?
- Take 10 minutes to think **why?!**
  - Hint: How many steps per `push_back`?
  - Hint: How many steps in all program?
  - Just approximate

```
int n = 4;
Vector v(n);
for (int i = 0; i < n; ++i)
    v.set(i, i);

v.push_back(15);
v.push_back(17);
v.push_back(19);
v.print();
// 0 1 2 3 15 17 19

for (int i = 0; i < 1000000; ++i) {
    v.push_back(i);
} // takes tooooo much time! WHY
```

# Approximately, How many steps

- Assume array has length size, If we tried to estimate the number of steps
  - $\sim 5 \times \text{size} + 7$
  - If size is  $10^6$ , this is around 5 millions.
  - In other words, it takes **linear number** of steps to be finished

- For simplicity, let's assume it takes size steps
  - E.g. we dropped constants

```
void push_back(int value) {  
    int *arr2 = new int[size + 1]; // size+1 steps  
    for (int i = 0; i < size; ++i) // 3size+1 steps  
        arr2[i] = arr[i];  
    arr2[size++] = value; // 2 steps  
    swap(arr, arr2); // 3 steps  
    delete[] arr2; // size steps  
    // Total: 5size + 7  
}
```

# Approximately, What is Overall number of steps?

- Now this loop is iterating  $n = 10^6$  step
  - We call push back, where size is increasing in each step, 1, 2, 3, .... n
  - So steps per a call are  $1 + 2 + 3 + \dots + n \Rightarrow \sim n * (n+1) / 2 = \mathbf{1/2 (n^2 + n)}$  steps
  - For simplicity, let's keep only the largest factor here  $\Rightarrow n^2$ 
    - Drop constant  $1/2$  and the smaller factor n
  - In other words, it takes **quadratic number** of steps for loop/body to be finished
  - So for  $n=10^6$ , we take  $\sim 10^{12}$  (multiplied by some factor, e.g. 5)
- Now, you know, **mathematically**, why this code takes too much time!

```
for (int i = 0; i < 1000000; ++i) {  
    v.push back(i);  
}
```

# Your turn

- In the next video, we will present a simple idea that speed the code
- Take 15 minutes to guess the trick
- Also, consider today analysis one good reason why we need to study the data structures NOT just use them!

*“Acquire knowledge and impart it to the people.”*

*“Seek knowledge from the Cradle to the Grave.”*