AMERICAN UNIVERSITY OF ARMENIA

College of Science and Engineering

CS 120 Introduction to Object-Oriented Programming

MIDTERM EXAM

Date / Time:

Friday, March 17 2017 at 17:30

Duration: Attention: 2 hours

ANY TYPE OF COMMUNICATION IS STRICTLY PROHIBITS
Write down your section, name and ID# at the top of all used pages

Participation:

Problem 1: Consider below a C++ function float kahan(float num1, float num2, float& compensation) that implements the Kahan Summation Algorithm for high-precision compensated summation of two float arguments float num1 and float num2:

float kahan(float num1, float num2, float &compensation)

float result;
num2 -= compensation;
result = num1 + num2;
compensation = (result - num1) - num2;
return result;

Using this function, write a C++ function *float pi(int n)* that computes the value π by the following

 $\pi = 16 \sum_{k=0}^{n} \frac{(-1)^{k}}{(2k+1)5^{2k+1}} - 4 \sum_{k=0}^{n} \frac{(-1)^{k}}{(2k+1)239^{2k+1}} = \left(\frac{16}{1*5} - \frac{4}{1*239}\right) - \left(\frac{16}{3*5^{3}} - \frac{4}{3*239^{3}}\right) + \left(\frac{16}{5*5^{5}} - \frac{4}{5*239^{5}}\right) - \cdots$ $+ \ln \text{clade} \qquad \text{with } h'$

The initial value of float compensation is 0.0.

[lead pi (int n) {

 float result; float compansation = 0.0;

 for (int c=0; C <= n, (++) h

 float and 1 = (mult pozo (-1, c)) / (2 c + 1) * madh pows (5, 2(+1);

 float and 2 = (mult powe (-1, c)) / (3 c + 1) * mult pows (259, 2 c + 1);

 float step = 16 ever * and 1 - 4 * and 2;

 result = kahan (usult, step, compansation)

When the backside if needed

Problem 1 of 4

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Problem 2: Write a Java method public static double[] expReg(double[] data) that takes as its argument an array of data points double[] data, and returns a two-element array - the first element being the exponent of an exponential regression and the second element being the amplitude. The exponential regression approximates the data points by a formula

where the exponent m and the amplitude a are computed as

 $m = \frac{\overline{xy - \overline{x}} \, \overline{y}}{\overline{z} - \overline{z}}, a = \overline{y} - m \, \overline{x}$

x y = mell (0) (04)

Here \bar{x} is the mean of the x coordinates, \bar{y} is the mean of the natural logarithm of y coordinates, \bar{x}^2 is the mean of the squares of the x coordinates, and \overline{xy} is the mean of the products of the x and natural logarithm of y coordinates. Use the element indices of the array double[] data as x coordinates and the element values as y coordinates. For natural logarithm, use the method double Math.log().

Both result elements are zeros, if at least one data element is non-positive.

public valic double [] expreg (double [] dota) {

double enploy double [] [] oursurer = mune double [] [double longth]; fore (int x = 0, x & double longth, x+1){ for (int x=0, x & double length, x++) 2 for (1ht X=0, x = muth. log (dida [x]); if negative?

double by = Muth. log (ny);

for (nxy=0,4chy)fill double m = ((x x ny)) / (x

double xM = [(1x x ny)12) f (x x ny)) / x x - oy x dy

louble m = ((1x x ny)12) f (x x ny)) / x x - oy x dy double a = dy - m & dx x; answer [0] [= m; 4 aurmen [1][0 = a

ansuler; Use the backside, if needed

Problem 2 of 4

Problem 3: Write a Java function public static boolean isInside(double[][] vertex, double x, double y) that takes as its argument a 2-by-n array of a convex polygon's vertex coordinates double[][] vertex - the x coordinates in the first row and y coordinates in the second row, and double x and double y coordinates of a point. It checks, if the point is inside the polygon.

Assume and use a method boolean to Left (double x1, double y1, double x2, double y2, double x0, double y0) that takes as its arguments coordinates of three points and returns true, if the third point (x0, y0)is in the left-hand side, when moving from the first point (x1, y1) to the second one (x2, y2); and false, if it 6 point

public static boolean is Inside bloable [][] very facte x obis is in the right-hand side. bool result: bool [] check Array = New bool [darketarray];
bool check, - New bool [Nouth I ow down town];

bool check, - New bool [Nouth I ow down the them I length (1))

if the contract of the down the them I length (1)

vielex [con] [al+1], [rowth][al+1],

y else 2 state = to Left (verdex [row][wl], verdex [d][wl], x, y)

check [w] = state;

1x Assume there is a function checking if all the value in the array are true or balse. If one of the value soil the same that are others it returns gabe.

result = is Some (check)

Use the backside, if needed

Problem 3 of 4

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Problem 4: Write a Java method public static void magicOdd(int[][] square) that creates a magic square of
an odd size using the following algorithm: 1. The number I goes in the middle of the top row; V Clear (IH[][] orr) - cluss III
2. All numbers are then placed one column to the right and one row up from the previous number;
number in the bottom row (note the place of 2 instead of the shaded location); $\sqrt{}$
4. Whenever the next number placement is outside of the rightmost column, stay in the same row and place the number in the leftmost column (note the place of 3 instead of the shaded location); $\sqrt{}$
1 U at t-1-w the measurement

5. When encountering an already filled-in square, place the next number directly below the previous number;

6. When the next number position is outside both a row and a column, place the number directly beneath the previous number (note the place of 7 instead of the shaded location).

beneath the previous number (note the place of / instead of the snaded location).
public static void magic Odd (int [] [] square) { 8 1 6 8 3 5 7 3 1000 = 0 100 k, l, m 100 list cquare (again) 100 lis
first C = 0
else if (now = square length) { -cow = square length 1 else if (square [cow] [square [row] [firstc] &c) t
7000 = 1000 + 2; $6000 = 6000 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1$
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