AMERICAN UNIVERSITY OF ARMENIA

College of Science and Engineering

CS 120 Introduction to Object-Oriented Programming

MIDTERM EXAM

Date / Time:

Friday, March 17 2017 at 17:30

Duration:

2 hours

Attention:

ANY TYPE OF COMMUNICATION IS STRICTLY PROHIBITED Write down your section, name and ID# at the top of all used pages

Participation:

Problem 1: Consider below a C++ function float kahan(float num1, float num2, float& compensation) that implements the Kahan Summation Algorithm for high-precision compensated summation of two float arguments float num1 and float num2:

float kahan (float num1, float num2, float &compensation)

float result; num2 -= compensation; result = num1 + num2; compensation = (result - num1) - num2; return result;

Using this function, write a C++ function float pi(int n) that computes the value π by the following

formula:

$$\pi = 16 \sum_{k=0}^{n} \frac{(-1)^{k}}{(2k+1)5^{2k+1}} - 4 \sum_{k=0}^{n} \frac{(-1)^{k}}{(2k+1)239^{2k+1}} = \underbrace{\left(\frac{16}{1*5} - \frac{4}{1*239}\right)}_{1} + \underbrace{\left(\frac{16}{3*5^{3}} - \frac{4}{3*239^{3}}\right)}_{3*239^{3}} + \underbrace{\left(\frac{16}{5*5^{3}} - \frac{4}{5*239^{3}}\right)}_{1} + \underbrace{\left(\frac{1$$

The initial value of float compensation is 0.0.

Sloat pilintal of float sam = 0;

for(inf i=0; i.c=n; i++) € Ke float step = 169/1.01

float step = 16.0 f (. D. f. pow (5.0 f, i)

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Problem 1 of 4

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Sloat sum = 0

fl

for(int i=1; i <= n; i+=2) {

float stap & part1=(16/(i.pow(s,i)));

Sloat port2 = (16/(i.pow(239,i)));

Sloat stop = Kahan (port1, part2, 0.0f);

Sum = Kahan (sum, stop, 0,0f);

J

7 Itarn Sum;

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Problem 2: Write a Java method public static double[] expReg(double[] data) that takes as its argument an array of data points double[] data, and returns a two-element array – the first element being the exponent of an exponential regression and the second element being the amplitude. The exponential regression approximates the data points by a formula

where the exponent m and the amplitude \hat{a} are computed as

 $m = \frac{\overline{xy} - \overline{x}\,\overline{y}}{\overline{x^2} - \overline{z}^2}, a = \overline{y} - m\,\overline{x}$

Here \overline{x} is the mean of the x coordinates, \overline{y} is the mean of the natural logarithm of y coordinates, $\overline{x^2}$ is the mean of the squares of the x coordinates, and \overline{xy} is the mean of the products of the x and natural logarithm of y coordinates. Use the element indices of the array double [] data as x coordinates and the element values as y coordinates. For natural logarithm, use the method double Math.log().

Both result elements are zeros, if at least one data element is non-positive.

Find m, X

$$\overline{g} = Moth. log(array Saluk)$$
 \overline{X}
 $\overline{X} = \frac{sum(ords)}{n}$

Sloat expreg (double [] dat) { float xsam = 0 j &s floor ysam = 0; floor f gmlar = 0; float xsqsum= 01 Sloat Xgmlan =0 int n = data-length; forlint i=0; icn; it+) of Sloat dat = data[i]; if (dat < 0) f Use the backside, if needed Tital new Sloat (] {0,0};

Problem 2 of 4

```
Xsum +=i;
  your + tadat;
  ymean += Math. log(dat);
  xsqsum + = Math. powldat, 2);
  xymean = i * Math log(dat);
X Sum = X Sum / (Sloat)n;
gmdan = gmean/(Sloot)nj
x sq sum = x sq sum/(Sloat/n;
xgmlan = xgmlan/Woatln;
float m = (xgmean - (gmean . xsum))/(xsqsum - Math.pou (xsum,2));
Sloat a = ymean - M*xsum;
  ulturn neu Sloat [] {m, a};
```

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Problem 3: Write a Java function public static boolean is Inside (double []] vertex, double x, double y) that takes as its argument a 2-by-n array of a convex polygon's vertex coordinates double [][] vertex — the x coordinates in the first row and y coordinates in the second row, and double x and double y coordinates of a point. It checks, if the point is inside the polygon. Assume and use a method boolean to Left (double x1, double y1, double x2, double y2, double x0, double y0) that takes as its arguments coordinates of three points and returns true, if the third point (x0, y0) is in the left-hand side, when moving from the first point (x1, y1) to the second one (x2, y2); and false, if it is in the right-hand side.
is Inside (double [] [] vertex, doublex, doubley) {
Flod initi
bool. initial value = true;
for (int i=0, i = fertex.length-vertex[0].length; i+=2){
isting the second of the secon
initial and -
Value (Co. 2) C
int if (i==0) { Ligitial value = ibleff (Vertex[i][d], Vertex[i][]);
int it(c==0) { Linitial value collected (value collected co
Continul;
iffis bool Isleft = isleft (vertex [i][0]) vertex[i][0])
Vartex[i+1][0], vartex[i+1][1],x',y);
Virtex[i+1][0], virtex[i+1][1], x', y); if ([slift]=initial value) \mathcal{E} Problem 3 of 4
Use the backside, if needed 71 tara false;

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forlint i= Vertexes[0]. length-1; i7/1;i-22){ if (i == Vertexes(O]. length-1) & inItial value = isleft (Ventexes[i]); vertex es[i]] vertexes[4][0], vertexa-[i-1][1]; x,9); continuli bool Isballot = islest (vertexes CiJCOJ, vertex es [i]COJ, vertexes[i-1][0], vertexes[i-1][1], x,y/i if (Isleft = initialvalue) & Utarn Salse; elfarn tralj

Problem 4: Write a Java method public static void magicOdd(int[][] square) that creates a magic square of an odd size using the following algorithm:

1. The number I goes in the middle of the top row;

2. All numbers are then placed one *column to the right* and *one row up* from the previous number;

3. Whenever the next number placement is above the top row, stay in the same column and place the number in the bottom row (note the place of 2 instead of the shaded location);

4. Whenever the next number placement is outside of the rightmost column, stay in the same row and place the number in the leftmost column (note the place of 3 instead of the shaded location);

5. When encountering an already filled-in square, place the next number directly below the previous number;

6. When the next number position is outside both a row and a column, place the number directly beneath the previous number (note the place of 7 instead of the shaded location).

beneath the previous number (note the place of 7 instead of the shaded location).	
magic Odd (int[][] square/q	-
int n=square[0]. lingth;	_
int positions= (n/2)(+1).	
int positions=0;	
for(int i=0; i <pow(n,2); i++){<="" td=""><td>-</td></pow(n,2);>	-
- for (int i=1; i = poor (n, 2) j i + +) {	
if (positionx in) {	
position x = 0;	
3	
if(positiong<0)	
positiong = N-1;	
3	
square[positiony][positionx]=i;	
positiony-=1; positionx +=1; overwrite?	
positionx += 1; overwrite?	Pro

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Problem 4 of 4

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