Memory Management Background: 1. Computer System Review 2. Address Binding & Linking ECE595 Sep 27 Y. Charlie Hu

"Things related to memory" -you have learned/heard so far



- · What is a processor
- What are registers
- What is memory
- How's memory organized
- What's a heap?
- What's a stack?
- Globals, locals, etc.
- PC, SP
- All of above deal with logical memory!
- Hardware memory
 What's a cache and how's it organized
 Physical memory a whole new can of worms!

1

Warning! You May Be Bored



- This material may be redundant if
- You've already had it (but may have forgotten)
- You already hacked and found it
- Your first language was assembly
- Feel free to ...
 - I won't be offended
 - You'll still be held responsible for the material

• C, Unix, Uniprocessors, No Threads

2

Warning: Approximate Truth



- Some details for general info
- Most details ignored entirely
- Goals
 - Simplicity
- Coverage

What does a Processor Do?



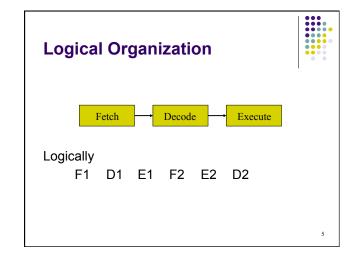
while (1)

fetch (get instruction)

decode (understand instruction)

execute

Execute: load, store, test, math, branch



Processor Operations



Logically

F1 D1 E1 F2 D2 E2

Pipeline

F1 D1 E1

F2 D2 E2

F3 D3 E3

What is the condition for a smooth pipelining? What can happen to the E part of Load inst: LD R0, _Y

What Is Memory (Address Space)



- "Slots" that hold values
- Slots are "numbered"
- Numbers are called addresses
- Two operations read or write
 - e.g., LD R1, _X
- What can you put in memory?
 - Anything. No "intrinsic" meaning

What is Cache?



- Another kind of "memory", physically
- Closer to the processor
- More expensive, smaller, faster
- Operation is logically transparent
 - No naming/access by
 - program, compiler, linker, loader
 - OS?
 - CPU?

CPU Where do we hope requests get satisfied? e.g. LD R1, _X L2 Main Memory

(Hardware) Memory Hierarchy

What Are Registers?



{

}

int i;

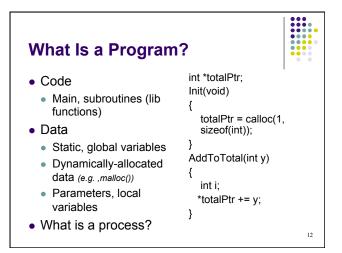
- Places to hold information
- Built into the processor
- "Named" specially

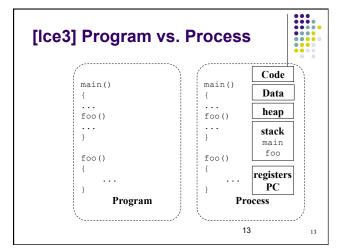
Why?

- Need a place to put operands / temp values • e = (a+b) * (c+d) * (a-f)
- · Highest level of memory hierarchy
- Register allocation problem NP-complete
 - · who does it?

What Is a Program?

```
int *totalPtr;
Init(void)
  totalPtr = calloc(1, sizeof(int));
AddToTotal(int y)
  *totalPtr += y;
```





Everything Becomes Memory



- Various ranges of memory (addr space) are used for different purposes
 - Text/Code (program instructions)
 - Data (global variables)
 - Stack (local variables, parameters, etc)
 - Heap (dynamically allocated memory)

14

What Is a Stack?



- Data structure that supports push/pop
- Uses?
 - Anything w/ LIFO (last-in first-out behavior)
 - · Only care about recent behavior
- Example?
 - DFS
 - Procedure calls!

Procedure calls



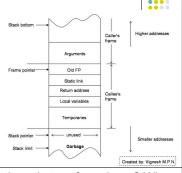
- Incoming parameters from caller
 - Don't even know who caller is
- · Local variables survive only when in use
- Temporary variables (a+b) * (c+d)

```
void Loop(int N)
{
    int a,b,c,d,e,f,g;
    ...
    g = (a+b*(c+d);
}
```

16

Stack Frames

- Frame = info for one procedure call
- Incoming parameters
- Return address for calle
- New local variables
- New temporary variable
- · Size of frame



• Where are the instructions that perform these? Who generate them if you program in C?

Stack Is Just Memory



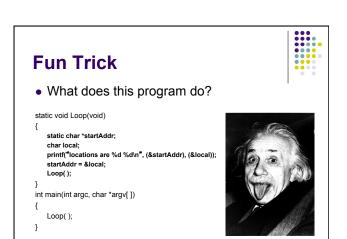
- Defined, used by convention
 - Agreement among OS, compiler, programmer
- How does OS manage stack (and heap)?
 - Allocate chunk of memory
 - Have pointer into chunk
- Problems?
 - Must know maximum size of stack?
- How to stay efficient despite uncertainty?

18

What Does Memory Look Like?



- Logical memory
 - · Code+data, stack, heap
 - Which ones grow?
 - How do you give them the most flexibility
- Physical memory?
 - · Another can of worms, entirely
- We will move on to Memory Management



Fun Trick



- Recursive function
- One static variable, one local variable
- Print difference between static variable and address of local variable
- Store address of local into static, recurs
- What does this tell you?
 - Address of local ?
 - · Address of static?
 - Address of main/Loop ?

```
A gap among Architecture,
Compiler and OS courses

main.c math.c math.o math.o
```

```
Example

Main.c:

main()
{
  static float x, val;
  extern float sin();
  extern printf(), scanf()
  printf("Type number: ");
  scanf("%f", &x);
  val = sin(x);
  printf("Sine is %f", val);
}

Math.c:

float sin(float x)
{
  static float temp1, temp2,
  result;
  - Calculate Sine -
  return result
}
```

Example (cont)

- Main.c uses externally defined sin() and C library function calls
 - printf()
 - scanf()
- How does this program get compiled and linked?

Tasks of a Linker



- Read in object files produced by the compiler
- Produce a self-sufficient object file (a.out)
 - Involves
 - segment relocation
 - address translation

25

Back to the Example



- Do the main/sin example with the following segment sizes
 - Main: code 420, data 42Math: code 1600, data 12
 - Library: code 1230, data 148
- Output: code 3250, data 202
- In reality segment starts on a page (4 Kbytes) boundary

26

Memory Layout – Division of Responsibility



- Compiler: generates object file
 - Information is incomplete
 - Each file may refer to symbols defined in other files
- Linker: puts everything together
 - Creates one object file that is complete
 - No references outside this file (usually)

Division of Responsibility (cont)



- OS (next major topic of this class)
 - Allow several different processes to <u>share</u> physical memory
 - Provide ways of <u>dynamically allocating more</u> physical memory

28

What could the compiler not do?



- Compiler does not know final memory layout
- It assumes everything in .o starts at address zero
- For each .o file, compiler puts information in the <u>symbol table</u> to tell the linker how to rearrange outside references safely/efficiently
 - For exported functions, absolute jumps, etc
- Linker needs to rearrange segments
- What makes rearrangement tricky?
 - Addresses!

29

What couldn't the compiler do? (cont)



- Compiler does not know all the references
 - e.g. addresses of functions / variables defined in other files
 - Where it does not know, it just puts a zero, and leaves a comment (<u>relocation info</u>) for the linker to fix things up
- These are called *cross references*

30

Components of Object File



- Header
- Two segments
 - Code segment and data segment
 - OS adds empty heap/stack segment while loading
- · Size and address of each segment
 - Address of a segment is the address where the segment begins

Components of Object File (cont)

- Symbol table
 - Information about stuff defined in this module
 - Used for getting from the name of a thing (subroutine/variable) to the thing itself
- Relocation information
 - Information about addresses in this module linker should fix
 - External references (e.g. lib call)
 - Internal references (e.g. absolute jumps)
- Additional information for debugger
 - . e.g. at breakpoint, "list"

32

Components of Object File (cont)



• Type "man 5 a.out" on UNIX for more information on UNIX object files

33

Linker functionality



34

- Three functions of a linker
 - Collect all the pieces of a program
 - Figure out new memory organization
 - Combine like segments
 - Does the ordering matter?
 - Touch up addresses
- The result is a runnable object file (e.g. a.out)

Linker – a closer look



 Linker can shuffle segments around at will, but cannot rearrange information within a segment

Recap: Tasks of a Linker



- Read in object files produced by the compiler
- Produce a self-sufficient object file (a.out)
- Implementation:
 - How many times does linker need to scan the object files?

36

Linker requires at least two passes



- Pass 1: decide how to arrange memory
- Pass 2: address touch-up

37

Pass 1 – Segment Relocation



- Pass 1 assigns input segment locations to fill up output segments
 - Read and adjust symbol table information
 - Read relocation info to see what additional stuff from libraries is required

38

Symbol Table



- Symbol table info:
 - Segments: name, size, old location, new location
 - Symbols: name, input segment containing it, offset within the segment

Pass 2 - Address translation



- In pass 2, linker reads segment and relocation information from files, fixes up addresses, and writes a new object file
- Relocation information is crucial for this part

40

Relocation Information



- Contains function/variable address and offest values to be relocated
- Examples of how to relocate:
 - "Place final address of symbol here"
 - LDW R1, _Y
 - Y is external
 - "Add difference between the final and original address of segment to the contents of this location"
 - J X
 - _X is in the original segment
 - "Add final address of symbol to contents of this location"
 - C reference of the form x = y.q;
 - y is external struct with {int p; int q}

41

Putting It Together



- Pass 1:
 - Read symbol table, relocation table
 - Rearrange segments, adjust symbol table
- Pass 2:
 - · Read segments and relocation information
 - Touch-up addresses
 - Write new object file

42

Course schedule (tentative)



Week 1	Intro / OS components	
Week 2	Processes	Lab 1
Week 3-4	Process Synchronization	Lab 2
Week 5-6	CPU scheduling, Deadlock, MM background	Lab 2
Week 7	Thread, Midterm	Lab 3
Week 8	Fall break, Memory management	Lab 3
Week 9-11	Memory Management	Lab 3
Week 12-15	File system	Lab 4
Week 16	Distributed file sys. / Course review	
Dec	Final exam	

Reading assignment



• All the prerequisite for ECE595