Obs Given a certificat clecher program T(v,C) and a polynomial time board f(iwi), there is a Problem circuit that comptly C(w,C) that's polynomial in 5720 relative to lwl. [duplicate a ore clock-cycle computation f(iwi) times].

reduction mps w to T(w) (other t)

-so it knows w and can wild a circuit to Now T(w,c) f(m)

for f(iw) the clack cycles:

Then show well iff T(w) & 3- TAT what we see if

iff JC T(w,c) output yes so

iff wfl because LENP

Divide and Conquer Algorithms -typically recursive - Livide original problem into 1 or more smaller problem of the same form - use recursion to conquer all smaller problems - recombing angues to arriver original publican

example sorting. A [i.n] a integers desire to regardance that A[i] = A[i] = --= A[in] but nointain sque collection in A[1-n] ~ Divide = ACAT ACAME = = ACM (linear in h) Recombin analyse rutine using a recurrence recursive calls T(n) = 2T(1) + cn merge $T(n) = c' n \lg n$

Dynamic Programming

To Sivide-and-conquer optimization

with memoization.

make a set of chains to that the result optimizes

some scoring function