

Exam 2

Mean

Std Dev

Whole Class

39!

18

Attend at least one of 2 times
we took attendance

45

Given a graph $G=(V, E)$

V is a set of vertices/nodes

$E \subseteq V \times V$ a set of edges (p, q) for $p \in V$ $q \in V$



a clique of size k in G is a set of k vertices $C \subseteq V$
s.t. every edge inside C is present, i.e. $C \times C \subseteq E_G$

Given a graph G and a size k , is there a size k clique in G ?

k -Clique = $\{ \langle G, k \rangle \mid G \text{ contains a size } k \text{ clique} \}$

— no better method known than "try all k -tuples of vertices" $(n^k \text{ tuples})$

$\text{CLIQUE-VAL} = \{ \langle G, k, C \rangle \mid C \text{ is a size } k \text{ clique in } G \} \in P$
(fact polynomial)

Observation: "Yes" and "No" are asymmetric here.

I can convince you that " G, k " \in k -Clique by
showing you a size k clique $C \subseteq V$ (and you check
Clique-eval)

But there is no clear way to show you " G, k " \notin k -Clique
without exponential effort (search all k -tuples)

Likewise I can certify that " φ " \in FORMULA-SAT
by showing truth assignment p s.t. $\varphi, p \in \text{FORMULA-EVAL}$
Scalable check
i.e. $\in P$