

Mathematical Induction

- for proving properties of the natural numbers $\forall n \geq c \ P(n)$ where $P(n)$ means "the property holds of n "

$$\begin{array}{r} 1 + \dots + n \\ n + \dots + 1 \\ \hline \underbrace{(n+1) + \dots + (n+1)}_{n \text{ times}} \\ n(n+1) \end{array}$$

$$\forall n \ n \geq c \rightarrow P(n)$$

$$\forall k \ \text{red}(k) \rightarrow \text{red}(k+1) \Rightarrow \forall n \geq c \ \text{red}(n)$$

$\text{red}(c)$

First Principle of Mathematical Induction

$$P(c) \wedge \left[\forall k \geq c \ P(k) \rightarrow P(k+1) \right] \rightarrow \forall n \geq c \ P(n)$$

base case

e.g. induction hypothesis (IH)

$$P(n) \equiv \left[\sum_{i=1}^n i = \frac{n(n+1)}{2} \right]$$

Show $\forall n \geq 0 \ P(n)$ where $P(n) \equiv \sum_{i=1}^n i = \frac{n(n+1)}{2}$

by induction,

Base case has $n=0$, show $(P(0), \text{i.e.}) \sum_{i=1}^0 i = \frac{0(0+1)}{2}$
 $0=0 \quad \checkmark$

Consider arbitrary $k \geq 0$ and show $P(k) \rightarrow P(k+1)$ using deduction method

Assume $P(k)$, i.e. $\sum_{i=1}^k i = \frac{k(k+1)}{2}$
induction hypothesis

and show $P(k+1)$, i.e. $\sum_{i=1}^{k+1} i = \frac{(k+1)(k+2)}{2}$


$$\sum_{i=1}^{k+1} i = \sum_{i=1}^k i + k+1 = \frac{k(k+1)}{2} + k+1 \quad \text{by IH} = \frac{(k+1)(k+2)}{2}$$

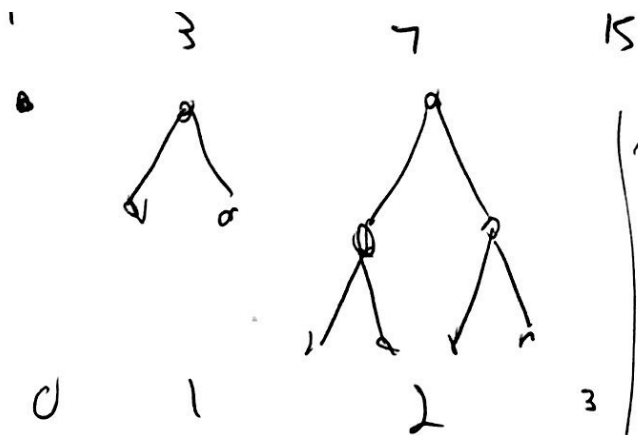
Prove $\forall n \geq 1, n < 2^n$
 $\forall n \geq 1, P(n)$

$$P(k) \equiv k < 2^k$$

Proof

$1 < 2^1 = 2$ so property holds at 1.

Suppose for induction $k < 2^k$ then $k+1 < 2^k + 1 \leq 2^k + 2^k = 2^{k+1}$
IH 



There are $\frac{2^{n+1}-1}{2}$ nodes in a complete binary tree of depth n .

Show by induction.

There are $2^{0+1}-1 = 1$ node in a depth 0 tree, for base case.

Consider a.h. $k \geq 0$ and suppose for induction there are $2^{k+1}-1$ nodes in a depth k binary tree. Then a depth $k+1$ binary tree has $2(2^{k+1}-1) + 1$ nodes, as desired \square

$$2^{k+1}-1$$

