Oubits:	/0>, 1>	instead of 45, (->	> for box is of 12	
Errors:	10>-b 11 14>-b 10>		Max	
	10> -> 10 1> -> -		nor	
	10>,14>-	L U107, U112 M	nitary continuous errors	
	Error Correction			
•	itition code 1970 Mariner Prob	e going to Mars -	greyscale pictura	
		rey -> 100 (
	A bits -> 12	bits, convert 1 emor		1 Wins
Hamm	ing Code			
	4 bits - > 7 #2 - Field w	bits correctury 2 elements	1-bit errors	
	addition mod			
		D * D= D D * 1= (* 0=0) * (=0		
	Fin - veetar sp.	ale over #2 relenents	(x_1, x_2, \dots, x_n) $x_i \in$	= F ₂
		(0,0),(1,0,		
codes	: A binary,	linear code C of le subspace of F2 ⁿ o	Pereth in and rank K, In, K1, of dimension k: 2k	<i>I</i> 5

To expense (Q,Q) , $(1,Q)(Q,Q)$, $(1,1)$ $E=7$ (Q,Q) , $(1,1)^3$ - Answ subspace of 15^2 2 vectors $1=2$ expenses $1=1$ work vector $(1,1)=1$ work $1=1$ wor						U	υ,	, u	V2	ł E	2			ω, ·	ŧω	2 t	e e																	
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Claim: Hamming Code can platest 2 bit error \$1 correct 1 bit error in length 7 word	A [<i>(</i>				0				1	1 1)	1			1				ı										
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