

数值算法 Homework 06

Due: Oct. 29, 2024

姓名: 雍崔扬

学号: 21307140051

Problem 01

Suppose that a tall-skinny matrix $A \in \mathbb{R}^{m \times n}$ is upper bidiagonal (i.e., $a_{ij} \neq 0$ if and only if $i - j \in \{-1, 0\}$)

Design an algorithm based on Givens rotations to solve the ridge regression problem

$$\min \|Ax - b\|_2^2 + \lambda \|x\|_2^2$$

where λ is a given positive number.

- 观察: 之所以设置 $A \in \mathbb{R}^{m \times n}$, 可能是因为我们目前掌握的仅仅是实数域上的 Givens 变换.

Solution:

岭回归问题, 即带 l_2 惩罚项的最小二乘问题, 本质上是标准的最小二乘问题:

$$\begin{aligned} & \min_x \|Ax - b\|_2^2 + \lambda \|x\|_2^2 \\ & \Leftrightarrow \\ & \min_x \left\| \begin{bmatrix} A \\ \sqrt{\lambda} I_n \end{bmatrix} x - \begin{bmatrix} b \\ 0_n \end{bmatrix} \right\|_2^2 \\ & \Leftrightarrow \\ & \min_x \|\tilde{A}x - \tilde{b}\|_2^2 \\ \text{where } \tilde{A} = & \begin{bmatrix} A \\ \sqrt{\lambda} I_n \end{bmatrix} \in \mathbb{R}^{(m+n) \times n} \text{ and } \tilde{b} = \begin{bmatrix} b \\ 0_n \end{bmatrix} \in \mathbb{R}^{m+n} \end{aligned}$$

等价问题的法方程为:

$$\begin{aligned} \begin{bmatrix} A \\ \sqrt{\lambda} I_n \end{bmatrix}^H \begin{bmatrix} A \\ \sqrt{\lambda} I_n \end{bmatrix} x &= \begin{bmatrix} A \\ \sqrt{\lambda} I_n \end{bmatrix}^H \begin{bmatrix} b \\ 0_n \end{bmatrix} \\ \Leftrightarrow \\ (A^H A + \lambda I_n)x &= A^H b \end{aligned}$$

因此岭回归解 $x_{\text{Ridge}} := (A^H A + \lambda I_n)^{-1} A^H b$

Algorithm 1

考虑到 \tilde{A} 的特殊结构, 我们可以很方便地通过 Givens 变换得到其 QR 分解.

为直观地展示这一事实, 我们考虑 $\begin{cases} m = 4 \\ n = 3 \end{cases}$ 的例子:

$$\tilde{A} = \begin{bmatrix} A \\ \sqrt{\lambda}I_n \end{bmatrix} = \begin{bmatrix} * & * & * \\ * & * & * \\ 0 & 0 & 0 \\ \hline * & & * \\ & * & * \end{bmatrix}$$

$$\Rightarrow G_{1,5}\tilde{A} = \begin{bmatrix} * & * & * \\ * & * & * \\ * & & * \\ 0 & 0 & 0 \\ 0 & * & * \\ * & & * \end{bmatrix} \Rightarrow G_{2,6}G_{2,5}(G_{1,5}\tilde{A}) = \begin{bmatrix} * & * & * \\ * & * & * \\ * & & * \\ 0 & 0 & 0 \\ 0 & 0 & * \\ 0 & * & * \end{bmatrix} \Rightarrow G_{3,7}G_{3,6}G_{3,5}(G_{2,6}G_{2,5}G_{1,5}\tilde{A}) = \begin{bmatrix} * & * & * \\ * & * & * \\ * & & * \\ 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix}$$

这样我们就得到 \tilde{A} 的 QR 分解为:

$$\tilde{A} = \tilde{Q}\tilde{R} \text{ where } \begin{cases} \tilde{R} = G_{3,7}G_{3,6}G_{3,5}(G_{2,6}G_{2,5}G_{1,5}\tilde{A}) \in \mathbb{R}^{(m+n) \times n} \\ \tilde{Q} = G_{1,5}^T G_{2,5}^T G_{2,6}^T G_{3,5}^T G_{3,6}^T G_{3,7}^T \in \mathbb{R}^{(m+n) \times (m+n)} \end{cases}$$

我们取 $\begin{cases} Q = \tilde{Q}(1 : m+n, 1 : n) \in \mathbb{R}^{(m+n) \times n} \\ R = \tilde{R}(1 : n, 1 : n) \in \mathbb{R}^{n \times n} \end{cases}$ 便得到 \tilde{A} 的精简 QR 分解.

使用回代法求解上三角方程组 $Rx = Q^T\tilde{b}$ 即可得到 $\min_x \|\tilde{A}x - \tilde{b}\|_2^2$ 问题的解.

(1) Helper Functions

生成尺寸为 $m \times n$ ($m \geq n$) 的上双对角矩阵 A 的函数:

```
function A = Upper_Bidiagonal(m, n)
    % Generates an upper bidiagonal matrix of size m x n
    % where m >= n. The main diagonal and the upper diagonal
    % contain random elements, and the remaining elements are zeros.
    %
    % Inputs:
    %   m - number of rows
    %   n - number of columns
    %
    % Output:
    %   A - resulting m x n upper bidiagonal matrix

    % Check if m is greater than or equal to n
    if m < n
        error('m must be greater than or equal to n');
    end

    % Construct the upper bidiagonal matrix
    A = diag(rand(n, 1)) + diag(rand(n-1, 1), 1);

    % Add rows of zeros to make it m x n
    A = [A; zeros(m - n, n)];
end
```

Givens 变换的算法为:

```

function: [c, s] = Givens(a, b)
    if b = 0
        c = 1; s = 0
    else
        if |b| > |a|
            t =  $\frac{a}{b}$ ; s =  $\frac{1}{\sqrt{1+t^2}}$ ; c = st
        else
            t =  $\frac{b}{a}$ ; c =  $\frac{1}{\sqrt{1+t^2}}$ ; s = ct
        end
    end
end

```

其 Matlab 代码如下:

```

function [c, s] = Givens(a, b)
% Givens 旋转, 计算 cos 和 sin
if b == 0
    c = 1;
    s = 0;
else
    if abs(b) > abs(a)
        t = a / b;
        s = 1 / sqrt(1 + t^2);
        c = s * t;
    else
        t = b / a;
        c = 1 / sqrt(1 + t^2);
        s = c * t;
    end
end
end

```

回代法的 Matlab 代码如下:

```

function x = Backward_Sweep(U, y)
% 回代法求解 UX = y
n = length(y);
for i = n:-1:2
    y(i) = y(i) / U(i, i); % 对角线归一化
    y(1:i-1) = y(1:i-1) - y(i) * U(1:i-1, i); % 消去
end
y(1) = y(1) / U(1, 1); % 处理第一行
x = y; % 返回结果
end

```

(2) 核心算法

生成增广矩阵 $\tilde{A} = \begin{bmatrix} A \\ \sqrt{\lambda}I_n \end{bmatrix}$ 并使用 Givens 变换计算其精简 QR 分解的函数 `Augmented_System_Fast_QR`:

$A \in \mathbb{R}^{m \times n}$ is an upper bidiagonal matrix where $m \geq n$, $\lambda \in \mathbb{R}$ is a positive number
function: $[Q, R] = \text{Augmented_System_Fast_QR}(A, \lambda)$

$$\tilde{A} = \begin{bmatrix} A \\ \sqrt{\lambda} I_n \end{bmatrix}$$

$$\tilde{Q} = I_n$$

for $j = 1 : n - 1$
 for $i = 1 : j$
 $[c, s] = \text{Givens}(\tilde{A}(j, j), \tilde{A}(m + i, j))$
 $\tilde{A}([j, m + i], j : j + 1) = \begin{bmatrix} c & s \\ -s & c \end{bmatrix} \tilde{A}([j, m + i], j : j + 1)$
 $\tilde{Q}(1 : m + n, [j, m + i]) = \tilde{Q}(1 : m + n, [j, m + i]) \begin{bmatrix} c & s \\ -s & c \end{bmatrix}^T$
 end
 end

for $i = 1 : n$ (Handling the case of $j = n$)
 $[c, s] = \text{Givens}(\tilde{A}(n, n), \tilde{A}(m + i, n))$
 $\tilde{A}([n, m + i], n) = \begin{bmatrix} c & s \\ -s & c \end{bmatrix} \tilde{A}([n, m + i], n)$
 $\tilde{Q}(1 : m + n, [n, m + i]) = \tilde{Q}(1 : m + n, [n, m + i]) \begin{bmatrix} c & s \\ -s & c \end{bmatrix}^T$
 end

$Q = \tilde{Q}(1 : m + n, 1 : n)$
 $R = \tilde{A}(1 : n, 1 : n)$

end

其 Matlab 代码为:

```

function [Q, R] = Augmented_System_Fast_QR(A, lambda)
    % Augmented_System_Fast_QR performs QR decomposition using Givens rotations.
    %
    % Inputs:
    %   A      - An upper bidiagonal matrix of size m x n (m >= n).
    %   Lambda - A positive scalar value.
    %
    % Outputs:
    %   Q - Orthogonal matrix of size (m+n) x n resulting from the QR decomposition.
    %   R - Upper triangular matrix of size n x n.

[m, n] = size(A); % Get dimensions of matrix A

% Check if m is greater than or equal to n
if m < n
    error('Matrix A must have m >= n.');
end

% Construct the augmented matrix
% \tilde{A} = [A; sqrt(lambda) * I_n]
I_n = eye(n); % Identity matrix of size n
A_tilde = [A; sqrt(lambda) * I_n]; % Augmented matrix

% Initialize \tilde{Q} as the identity matrix of size (m + n) x (m + n)
Q_tilde = eye(m + n);

% Perform Givens rotations for j = 1 to n-1
for j = 1:n-1
    for i = 1:j
        % Givens rotation
        [c, s] = givens(A_tilde(j, j), A_tilde(m + i, j));
        A_tilde([j, m + i], j:j+1) = [c s; -s c] * A_tilde([j, m + i], j:j+1);
        Q_tilde(1:m+n, [j, m+i]) = Q_tilde(1:m+n, [j, m+i]) * [c s; -s c];
    end
end

```

```

    % Compute Givens rotation parameters
    [c, s] = Givens(A_tilde(j, j), A_tilde(m + i, j));

    % Apply the Givens rotation to the augmented matrix
    A_tilde([j, m + i], j:j + 1) = [c, s; -s, c] * A_tilde([j, m + i], j:j +
1);

    % Apply the Givens rotation to the orthogonal matrix \tilde{Q}
    Q_tilde(1:m + n, [j, m + i]) = Q_tilde(1:m + n, [j, m + i]) * [c, s; -s,
c]';
end
end

% Handle the case for j = n
for i = 1:n
    % Compute Givens rotation parameters
    [c, s] = Givens(A_tilde(n, n), A_tilde(m + i, n));

    % Apply the Givens rotation to the augmented matrix
    A_tilde([n, m + i], n) = [c, s; -s, c] * A_tilde([n, m + i], n);

    % Apply the Givens rotation to the orthogonal matrix \tilde{Q}
    Q_tilde(1:m + n, [n, m + i]) = Q_tilde(1:m + n, [n, m + i]) * [c, s; -s, c]';
end

% Extract Q from \tilde{Q}
Q = Q_tilde(1:m + n, 1:n);

% Extract R from the upper triangular portion of \tilde{A}
R = A_tilde(1:n, 1:n);
end

```

计算岭回归解的函数 `Solve_Upper_Bidiagonal_Ridge_Regression`:

```

function x = Solve_Upper_Bidiagonal_Ridge_Regression(A, b, lambda)
    % Solve_Upper_Bidiagonal_Ridge_Regression solves the ridge regression problem
    % for an upper bidiagonal matrix A using QR decomposition.
    %
    % Inputs:
    %   A      - An upper bidiagonal matrix of size m x n (m >= n).
    %   b      - A column vector of size m x 1 representing the dependent variable.
    %   Lambda - A positive scalar for regularization.
    %
    % Outputs:
    %   x      - The estimated coefficients for the ridge regression.

    [m, n] = size(A); % Get dimensions of matrix A

    % Check if m is greater than or equal to n
    if m < n
        error('Matrix A must have m >= n.');?>
        % Ensure the matrix A has more rows than columns
    end

    % Augment the response vector b with zeros to match the size of A
    b_tilde = [b; zeros(n, 1)];

    % Perform reduced QR decomposition using the Augmented System approach
    [Q, R] = Augmented_System_Fast_QR(A, lambda);

```

```

    % Solve the upper triangular system using backward substitution
    x = Backward_Sweep(R, Q' * b_tilde); % Calculate the coefficients

end

```

(3) 运行结果

我们将 $x_{\text{Ridge}} = (A^T A + \lambda I_n)^{-1} A^T b$ 作为精确解，
与 `Solve_Upper_Bidiagonal_Ridge_Regression` 得到的计算解进行对比。
函数调用如下：

```

rng(51);
m = 1000;
n = 500;
lambda = 1;      % Regularization parameter for ridge regression

% Generate an upper bidiagonal matrix A of size m x n
A = Upper_Bidiagonal(m, n);

% Generate random dependent variable vector b of size m x 1
b = rand(m, 1);

% Solve ridge regression using the custom function
x = Solve_Upper_Bidiagonal_Ridge_Regression(A, b, lambda);

% calculate the exact solution for comparison (using explicit inversion)
x_exact = (A' * A + lambda * eye(n, n)) \ (A' * b);

% Display the norm of the difference between x and x_exact (Frobenius norm)
disp("计算解与精确解之差的 Frobenius 范数:")
disp(norm(x - x_exact, "fro"));

```

运行结果：

```

计算解与精确解之差的 Frobenius 范数:
2.3420e-14

```

将 `lambda` 的值从 `1` 改到 `0.1`，再次运行得到：

```

计算解与精确解之差的 Frobenius 范数:
3.0359e-14

```

这验证了上述算法的正确性。

对比: Algorithm 2

回忆起等价问题的法方程为：

$$(A^T A + \lambda I_n)x = A^T b$$

注意到 A 作为一个实上双对角矩阵，能够保证 $A^T A + \lambda I_n$ 是一个对称正定三对角阵。
我们可以利用 Givens 变换方便地得到 $A^T A + \lambda I_n$ 的 QR 分解。

为直观地展示这一事实，我们考虑 $\begin{cases} m \geq n \\ n = 4 \end{cases}$ 的例子：

$$A^T A + \lambda I_n = \begin{bmatrix} * & * & & \\ * & * & * & \\ * & * & * & \\ * & * & * & \end{bmatrix}$$

$$G_{1,2}(A^T A + \lambda I_n) = \begin{bmatrix} * & * & * & \\ 0 & * & * & \\ * & * & * & \\ * & * & * & \end{bmatrix}$$

$$G_{2,3}[G_{1,2}(A^T A + \lambda I_n)] = \begin{bmatrix} * & * & * & \\ 0 & * & * & * \\ 0 & * & * & \\ * & * & * & \end{bmatrix}$$

$$G_{3,4}[G_{2,3}G_{1,2}(A^T A + \lambda I_n)] = \begin{bmatrix} * & * & * & \\ 0 & * & * & * \\ 0 & * & * & \\ 0 & * & * & \end{bmatrix}$$

$$A^T A + \lambda I_n = QR \text{ where } \begin{cases} R = G_{3,4}[G_{2,3}G_{1,2}(A^T A + \lambda I_n)] \\ Q = G_{1,2}^T G_{2,3}^T G_{3,4}^T \end{cases}$$

(1) 核心算法

求解法方程系数矩阵 $A^T A + \lambda I_n$ 的 QR 分解的函数 `Normal_Equation_Fast_QR`:

```
function [Q, R] = Normal_Equation_Fast_QR(A, lambda)
    % Normal_Equation_Fast_QR performs QR decomposition on the normal equation
    % matrix of a linearly constrained least squares problem using Givens rotations.
    %
    % Inputs:
    %   A      - An upper bidiagonal matrix of size m x n, where m >= n.
    %   Lambda - A positive scalar regularization parameter for ridge regression.
    %
    % Outputs:
    %   Q      - An orthogonal matrix of size n x n resulting from the QR decomposition.
    %   R      - An upper triangular matrix of size n x n, representing the R factor
    %           in the QR decomposition of the regularized normal equation matrix.

    % Get the dimensions of matrix A
    [~, n] = size(A);

    % Construct the regularized normal equation matrix for ridge regression
    % R starts as (A' * A + lambda * I), where I is the n x n identity matrix
    R = A' * A + lambda * eye(n);

    % Initialize Q as the identity matrix of size n x n
    Q = eye(n);

    % Apply Givens rotations to create an upper triangular matrix R
    % Iterate over each column up to n-1 to zero out sub-diagonal elements
    for j = 1:n-1
        % Compute the Givens rotation parameters (c, s) to zero out R(j+1, j)
        [c, s] = Givens(R(j, j), R(j+1, j));

        % Apply the Givens rotation to the submatrix R(j:j+1, j:j+1)
        % This operation zeros out the (j+1, j) entry of R
        R(j:j+1, j:j+1) = [c, s; -s, c] * R(j:j+1, j:j+1);

        % Update Q with the transpose of the Givens rotation to maintain orthogonality
        Q = Q * [c, s; -s, c];
    end
end
```

```

        Q(1:n, j:j+1) = Q(1:n, j:j+1) * [c, s; -s, c]';
    end
end

```

计算岭回归解的函数 `Solve_Upper_Bidiagonal_Ridge_Regression`:

```

function x = Solve_Upper_Bidiagonal_Ridge_Regression(A, b, lambda)
    % Solve_Upper_Bidiagonal_Ridge_Regression solves the ridge regression problem
    % for an upper bidiagonal matrix A using QR decomposition.
    %
    % Inputs:
    %   A      - An upper bidiagonal matrix of size m x n (m >= n).
    %   b      - A column vector of size m x 1 representing the dependent variable.
    %   lambda - A positive scalar for regularization.
    %
    % Outputs:
    %   x      - The estimated coefficients for the ridge regression.

[m, n] = size(A); % Get dimensions of matrix A

% Check if m is greater than or equal to n
if m < n
    error('Matrix A must have m >= n.');?>
    % Ensure the matrix A has more rows than
columns
end

% Perform reduced QR decomposition using the Augmented System approach
[Q, R] = Normal_Equation_Fast_QR(A, lambda);

% Solve the upper triangular system using backward substitution
x = Backward_Sweep(R, Q' * (A' * b)); % Calculate the coefficients

end

```

(2) 运行结果

我们将 $x_{\text{Ridge}} = (A^T A + \lambda I_n)^{-1} A^T b$ 作为精确解，
与 `Solve_Upper_Bidiagonal_Ridge_Regression` 得到的计算解进行对比。
函数调用:

```

rng(51);
m = 1000;
n = 500;
lambda = 1;      % Regularization parameter for ridge regression

% Generate an upper bidiagonal matrix A of size m x n
A = Upper_Bidiagonal(m, n);

% Generate random dependent variable vector b of size m x 1
b = rand(m, 1);

% Solve ridge regression using the custom function
x = Solve_Upper_Bidiagonal_Ridge_Regression(A, b, lambda);

% calculate the exact solution for comparison (using explicit inversion)
x_exact = (A' * A + lambda * eye(n, n)) \ (A' * b);

% Display the norm of the difference between x and x_exact (Frobenius norm)
disp("计算解与精确解之差的 Frobenius 范数:")
disp(norm(x - x_exact, "fro"));

```

运行结果:

计算解与精确解之差的 Frobenius 范数:

0.1680

将 `lambda` 的值从 1 改到 0.1, 再次运行得到:

计算解与精确解之差的 Frobenius 范数:

45.5396

这说明算法 2 的数值稳定性较差, 即法方程不适合直接进行数值求解.

Problem 02 (第一个算法有误)

Consider the linearly constrained least squares problem $\min\{\|Ax - b\|_2^2 : Cx = d\}$

In the lecture we use the QR factorization of C^H to eliminate the linear constraint $Cx = d$ and reduce the problem to a standard least squares one.

Design an algorithm based on Gaussian elimination to eliminate the linear constraint.

Solution:

考虑线性等式约束的最小二乘问题:

$$\min_{Cx=d} \|Ax - b\|_2^2 \text{ where } \begin{cases} A \in \mathbb{C}^{m \times n} \text{ and } \operatorname{rank}(A) = n \leq m \\ b \in \mathbb{C}^m \\ C \in \mathbb{C}^{p \times n} \text{ and } \operatorname{rank}(C) = p < n \\ d \in \mathbb{C}^p \end{cases}$$

设根据 Gauss 消元法得到 $C \in \mathbb{C}^{p \times n}$ 的 LU 分解为:

$$C = LU \text{ where } \begin{cases} L \in \mathbb{C}^{p \times p} \\ U = [U_1, U_2] \in \mathbb{C}^{p \times n} \\ U_1 \in \mathbb{C}^{p \times p} \text{ is a upper triangular matrix} \\ U_2 \in \mathbb{C}^{p \times (n-p)} \end{cases}$$

对应地将 A, x 分块为:

$$\begin{cases} A = [A_1, A_2] \in \mathbb{C}^{m \times n} \\ x = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} \in \mathbb{C}^n \end{cases} \text{ where } \begin{cases} A_1 \in \mathbb{C}^{m \times p} \\ A_2 \in \mathbb{C}^{m \times (n-p)} \\ x_1 \in \mathbb{C}^p \\ x_2 \in \mathbb{C}^{n-p} \end{cases}$$

则我们有:

$$\begin{aligned} Cx = LUx &= L[U_1, U_2] \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = L(U_1 x_1 + U_2 x_2) = d \\ &\Leftrightarrow \\ x_1 &= U_1^{-1}(L^{-1}d - U_2 x_2) \end{aligned}$$

于是我们有:

$$\begin{aligned}
\min_{Cx=d} \|Ax - b\|_2^2 &\Leftrightarrow \min_{Cx=d} \left\| [A_1, A_2] \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} - b \right\|_2^2 \\
&\Leftrightarrow \min_{Cx=d} \|A_1x_1 + A_2x_2 - b\|_2^2 \\
&\Leftrightarrow \min_{x_2 \in \mathbb{R}^{n-p}} \|A_1U_1^{-1}(L^{-1}d - U_2x_2) + A_2x_2 - b\|_2^2 \\
&\Leftrightarrow \min_{x_2 \in \mathbb{R}^{n-p}} \|(A_2 - A_1U_1^{-1}U_2)x_2 - (b - A_1U_1^{-1}L^{-1}d)\|_2^2 \\
&\Leftrightarrow \min_{x_2 \in \mathbb{R}^{n-p}} \|\tilde{A}x_2 - \tilde{b}\|_2^2
\end{aligned}$$

where $\begin{cases} \tilde{A} = A_2 - A_1U_1^{-1}U_2 \in \mathbb{C}^{m \times (n-p)} \\ \tilde{b} = b - A_1U_1^{-1}L^{-1}d \in \mathbb{C}^m \end{cases}$

(1) Helper Functions

生成线性等式约束的最小二乘问题的函数:

```

function [A, b, C, d] = Linearly_Constrained_Least_Squares(n, m, p)
    % Linearly_Constrained_Least_Squares creates a constrained least squares problem.
    %
    % Inputs:
    %   n - Number of variables (size of x).
    %   m - Number of observations (size of b).
    %   p - Number of constraints (size of d).
    %
    % Outputs:
    %   A - A full-rank matrix of size m x n.
    %   b - A vector of size m x 1.
    %   C - A rank-deficient matrix of size p x n.
    %   d - A vector of size p x 1.

    % Check input constraints
    if n > m
        error('Condition n <= m is violated: n must be less than or equal to m.');
    end
    if p >= n
        error('Condition p < n is violated: p must be less than n.');
    end

    % Generate a random full-rank matrix A (m x n)
    A = rand(m, n) + 1i * rand(m, n); % Adding complex components
    A = A * rand(n, n); % Ensure full rank by multiplying with a full-rank matrix

    % Check the rank of A
    if rank(A) ~= n
        error('Rank of A must be n.');
    end

    % Generate a random vector b (m x 1)
    b = rand(m, 1) + 1i * rand(m, 1); % Complex vector

    % Generate a random rank-deficient matrix C (p x n)
    % We can ensure C has rank p < n by making the last (n - p) columns zero
    C_temp = rand(p, p) + 1i * rand(p, p); % Create a random full-rank matrix
    C = [C_temp, zeros(p, n - p)]; % Append zero columns to make it p x n

    % Check the rank of C
    if rank(C) ~= p
        error('Rank of C must be p.');
    end

```

```
% Generate a random vector d (p x 1)
d = rand(p, 1) + 1i * rand(p, 1); % Complex vector
end
```

长方矩阵的 Gauss 消去法:

```
function: A = Gaussian_Elimination(A)
[m, n] = size(A)
for k = 1 : min{m - 1, n - 1}
    A(k + 1 : m, k) ← A(k + 1 : m, k) / A(k, k)
    A(k + 1 : m, k + 1 : n) ← A(k + 1 : m, k + 1 : n) - A(k + 1 : m, k)A(k, k + 1 : n)
end
end
```

长方矩阵 Gauss 消去法的 Matlab 代码如下:

```
function A = Gaussian_Elimination(A)
% Gaussian_Elimination performs Gaussian elimination on matrix A
%
% Inputs:
%   A - An m x n matrix to be transformed into an upper triangular form.
%
% Outputs:
%   A - The matrix after Gaussian elimination, with zeros below the main diagonal.

% Get dimensions of A
[m, n] = size(A);

% Perform Gaussian elimination
for k = 1:min(m-1, n-1)
    % scale the column below the diagonal in the k-th column
    A(k+1:m, k) = A(k+1:m, k) / A(k, k);

    % update the submatrix to eliminate entries in the k-th column
    A(k+1:m, k+1:n) = A(k+1:m, k+1:n) - A(k+1:m, k) * A(k, k+1:n);
end
end
```

前代法:

```
function y = Forward_Sweep(L, b)
% 前代法求解 Ly = b
n = length(b);
for i = 1:n-1
    b(i) = b(i) / L(i, i); % 对角线归一化
    b(i+1:n) = b(i+1:n) - b(i) * L(i+1:n, i); % 消去
end
b(n) = b(n) / L(n, n); % 处理最后一行
y = b; % 返回结果
end
```

回代法:

```

function X = Backward_Sweep(U, Y)
    % 回代法求解 UX = Y
    [n, ~] = size(Y);
    for i = n:-1:2
        Y(i,:) = Y(i,:) / U(i, i); % 对角线归一化
        Y(1:i-1,:) = Y(1:i-1,:) - U(1:i-1, i) * Y(i,:);
    end
    Y(1,:) = Y(1,:) / U(1, 1); % 处理第一行
    X = Y; % 返回结果
end

```

以下是基于复数域上的 Householder QR 分解的求解最小二乘问题的函数:

复数域上的 Householder 变换的计算算法已在 Homework 4 Problem 2 中给出:

```

function [v, beta] = Complex_Householder(x)
    % This function computes the Householder vector 'v' and scalar 'beta' for
    % a given complex vector 'x'. This transformation is used to create zeros
    % below the first element of 'x' by reflecting 'x' along a specific direction.

    n = length(x);
    x = x / norm(x, inf); % Normalize x by its infinity norm to avoid numerical issues

    % Copy all elements of 'x' except the first into 'v'
    v = zeros(n, 1);
    v(2:n) = x(2:n);

    % Compute sigma as the squared 2-norm of the elements of x starting from the second
    % element
    sigma = norm(x(2:n), 2)^2;

    % Check if sigma is near zero, which would mean 'x' is already close to a scalar
    % multiple of e_1
    if sigma < 1e-10
        beta = 0; % If sigma is close to zero, set beta to zero (no transformation
        % needed)
    else
        % Determine gamma to account for the argument of complex number x(1)
        if abs(x(1)) < 1e-10
            gamma = 1; % If x(1) is close to zero, set gamma to 1
        else
            gamma = x(1) / abs(x(1)); % otherwise, set gamma to x(1) divided by its
            % magnitude
        end

        % Compute alpha as the Euclidean norm of x, including x(1) and sigma
        alpha = sqrt(abs(x(1))^2 + sigma);

        % Compute the first element of 'v' to avoid numerical cancellation
        v(1) = -gamma * sigma / (abs(x(1)) + alpha);

        % Calculate 'beta', the scaling factor of the Householder transformation
        beta = 2 * abs(v(1))^2 / (abs(v(1))^2 + sigma);

        % Normalize the vector 'v' by v(1) to ensure that the first element is 1,
        % allowing for simplified storage and computation of the transformation
        v = v / v(1);
    end
end

```

复数域上的 Householder QR 算法已在 Homework 4 Problem 3 中给出:

```

function [Q, R] = Complex_Householder_QR(A)
[m, n] = size(A);
Q = eye(m); % Initialize Q as the identity matrix
R = A; % Initialize R as A

for k = 1:min(m-1, n)
    [v, beta] = Complex_Householder(R(k:m, k)); % Apply Complex Householder

    % Update R
    R(k:m, k:n) = R(k:m, k:n) - (beta * v) * (v' * R(k:m, k:n));

    % Update Q
    Q(1:m, k:m) = Q(1:m, k:m) - (Q(1:m, k:m) * v) * (beta * v');
end
end

```

计算得到 $A \in \mathbb{C}^{m \times n}$ 的 QR 分解 $A = QR$ 之后 (其中 $Q \in \mathbb{C}^{m \times m}$ 为酉矩阵, $R \in \mathbb{C}^{m \times n}$ 的上 $n \times n$ 分块 R_1 为上三角阵)

求解法方程 $A^H Ax = R^H Rx = A^H b$ 就等价于求解 $\begin{cases} R^H y = A^H b \\ Rx = y \end{cases}$ (分别由前代法和回代法求解)

或者也可考虑精简 QR 分解 $A = Q_1 R_1$ (其中 $Q_1 \in \mathbb{C}^{m \times n}$ 由 Q 的前 n 列构成)

则求解法方程 $A^H Ax = R_1^H R_1 x = R_1^H Q_1^H b = A^H b$ 就等价于求解 $R_1 x = Q_1^H b$ (由回代法求解)

```

function x = Householder_Solution(A, b)
[m, n] = size(A);

% Step 1: Compute the QR decomposition of A using Householder reflections
[Q, R] = Complex_Householder_QR(A);

% Step 2: Solve the system Rx = Q' * b using backward substitution
x = Backward_Sweep(R(1:n, 1:n), Q(1:m, 1:n)' * b);
end

```

(2) 核心算法

基于 Gauss 消去法的解法:

- ① 计算 $C \in \mathbb{C}^{p \times n}$ ($p < n$) 的 LU 分解
 $C = LU$ where $\begin{cases} L \in \mathbb{C}^{p \times p} \\ U = [U_1, U_2] \in \mathbb{C}^{p \times n} \\ U_1 \in \mathbb{C}^{p \times p} \text{ is a upper triangular matrix} \\ U_2 \in \mathbb{C}^{p \times (n-p)} \end{cases}$
- ② 计算 $\begin{cases} \tilde{A} = A_2 - A_1 U_1^{-1} U_2 \in \mathbb{C}^{m \times (n-p)} \\ \tilde{b} = b - A_1 U_1^{-1} L^{-1} d \in \mathbb{C}^m \end{cases}$
- ③ 计算 $x_2 = \text{Householder_Solution}(\tilde{A}, \tilde{b})$
(这里通过调用 Homework 5 Problem 6 的 `Householder_solution` 函数得到消元后的最小二乘问题的解)
- ④ 计算 $x_1 = U_1^{-1}(L^{-1}d - U_2 x_2)$, 得到 $x_{\text{Gauss}} = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$

```

function [x_Gauss, b_tilde, A_tilde] = Gaussian_Method(A, b, c, d)
% Gaussian_Method eliminates constrained variables in a linearly constrained
% least squares problem using Gaussian elimination.
%
% Inputs:

```

```

% A - The original matrix of size m x n, representing the coefficients of the
least squares problem.
% b - The original vector of size m x 1, representing the observed values.
% C - The constraint matrix of size p x n, representing the constraints on the
variables.
% d - The constraint vector of size p x 1, representing the values of the
constraints.
%
% Outputs:
% x_Gauss - The solution vector after eliminating variables.
% b_tilde - The modified vector after eliminating constrained variables.
% A_tilde - The modified matrix after eliminating constrained variables.

% Get dimensions of A and C
[m, n] = size(A); % m: number of observations, n: number of variables
[p, ~] = size(C); % p: number of constraints

% Step 1: Perform Gaussian elimination on C
C = Gaussian_Elimination(C); % Transform C into an upper triangular matrix

% Step 2: Extract L and U from the Gaussian elimination of C
L = eye(p) + triu(C(1:p, 1:p), -1); % Construct the lower triangular matrix L with
elimination factors
U = triu(C(1:p, 1:n)); % Extract the upper triangular matrix U

% Step 3: Compute the modified vector b_tilde based on the elimination
% d_tilde = U(1:p, 1:p) \ (L \ d);
d_tilde = Forward_Sweep(L, d);
d_tilde = Backward_Sweep(U(1:p, 1:p), d_tilde);
b_tilde = b - A(1:m, 1:p) * d_tilde; % Adjust b by removing the effect of the
constraints

% Step 4: Compute the modified matrix A_tilde to reflect the elimination
% U_tilde = U(1:p, 1:p) \ U(1:p, p+1:n);
U_tilde = Backward_Sweep(U(1:p, 1:p), U(1:p, p+1:n));
A_tilde = A(1:m, p+1:n) - A(1:m, 1:p) * U_tilde; % Adjust A to eliminate variables
according to constraints

% Step 5: Solve for x2 using the modified A_tilde and b_tilde
x2 = Householder_Solution(A_tilde, b_tilde); % Solve the least squares problem for
the remaining variables

% Step 6: Compute x1 from d_tilde and U_tilde
x1 = d_tilde - U_tilde * x2; % Calculate x1 based on the relationship with x2

% Combine x1 and x2 into the final solution vector
x_Gauss = [x1; x2]; % Concatenate x1 and x2 to form the complete solution vector
end

```

(3) 对比: Lagrange 乘子法

考虑线性等式约束的最小二乘问题:

$$\min_{Cx=d} \|Ax - b\|_2^2 \text{ where } \begin{cases} A \in \mathbb{C}^{m \times n} \text{ and } \operatorname{rank}(A) = n \leq m \\ b \in \mathbb{C}^m \\ C \in \mathbb{C}^{p \times n} \text{ and } \operatorname{rank}(C) = p < n \\ d \in \mathbb{C}^p \end{cases}$$

注意到目标函数 $f(x) = \|Ax - b\|_2^2$ 是关于 x 的凸函数, 而问题只有线性等式约束 $Cx = d$ 因此这是一个标准形式的凸优化问题, 其最优解即为 KKT 点.

定义其 Lagrange 函数 $L(x, \lambda)$ 为:

$$\begin{aligned} L(x, \lambda) &= f(x) - \lambda^H(Cx - d) \\ &= \|Ax - b\|_2^2 - \lambda^H(Cx - d) \\ \text{dom}\{L\} &= \mathbb{C}^n \times \mathbb{C}^p \end{aligned}$$

Lagrange 函数 $L(x, \lambda)$ 关于 x 的梯度为:

$$\begin{aligned} \nabla_x L(x, \lambda) &= \nabla_x \{\|Ax - b\|_2^2 - \lambda^H(Cx - d)\} \\ &= -A^H \cdot 2(b - Ax) - (\lambda^H C)^H \\ &= -2A^H b + 2A^H A x - C^H \lambda \end{aligned}$$

KKT 条件为:

$$\begin{cases} \nabla_x L(x, \lambda) = -2A^H b + 2A^H A x - C^H \lambda = 0_n & \textcircled{1} \\ Cx = d & \textcircled{2} \end{cases}$$

① 式左乘 $(A^H A)^{-1}$ 可得 $-2(A^H A)^{-1} A^H b + 2x - (A^H A)^{-1} C^H \lambda = 0_n$

于是有 $x = (A^H A)^{-1} A^H b + \frac{1}{2}(A^H A)^{-1} C^H \lambda$

代入 ② 式即得 $Cx = C(A^H A)^{-1} A^H b + \frac{1}{2}C(A^H A)^{-1} C^H \lambda = d$

解得 $\lambda_{\text{KKT}} = 2[C(A^H A)^{-1} C^H]^{-1}[d - C(A^H A)^{-1} A^H b]$

因此我们有:

$$\begin{aligned} x_{\text{Lagrange}} &= x_{\text{KKT}} \\ &= (A^H A)^{-1} A^H b + \frac{1}{2}(A^H A)^{-1} C^H \lambda_{\text{KKT}} \\ &= (A^H A)^{-1} A^H b + \frac{1}{2}(A^H A)^{-1} C^H \cdot 2[C(A^H A)^{-1} C^H]^{-1}[d - C(A^H A)^{-1} A^H b] \\ &= x_{\text{ls}} - (A^H A)^{-1} C^H [C(A^H A)^{-1} C^H]^{-1}(Cx_{\text{ls}} - d) \end{aligned}$$

其中 $x_{\text{ls}} := (A^H A)^{-1} A^H b$

我们将 x_{Lagrange} 视为问题的精确解, 与计算解 x_{Gauss} 进行比较.

以下是基于 Lagrange 乘子法的解法:

(简单起见, 我们直接使用 Matlab 内置的 \ 运算来求逆)

```
% Function to solve the constrained least squares problem using the Lagrangian method
function x_Lagrange = Lagrangian_Method(A, b, C, d)
    % Calculate the least squares solution without constraints
    x_ls = (A' * A) \ (A' * b); % Normal equation solution for least squares

    % Intermediate calculation to prepare for the constrained adjustment
    tmp = C * ((A' * A) \ C'); % Project the constraints into the least squares space

    % Adjust the least squares solution for the constraints
    x_Lagrange = x_ls - (A' * A) \ (C' * (tmp \ (C * x_ls - d)));
    % Apply the Lagrangian adjustment based on the constraints
end
```

(4) 运行结果

我们通过对基于 Gauss 消去法的解 x_{Gauss} 和基于 Lagrange 乘子法的解 x_{Lagrange} (视为精确解) 来检验我们结果的正确性.

函数调用:

```
rng(51);
m = 500; % Number of observations
n = 50; % Number of variables
p = 10; % Number of constraints
```

```

% Generate the linearly constrained least squares problem
[A, b, C, d] = Linearly_Constrained_Least_Squares(n, m, p);

% Apply Gaussian elimination method to eliminate constraints
[x_Gauss, b_tilde, A_tilde] = Gaussian_Method(A, b, C, d);

% Apply Lagrangian method to solve the same constrained problem
x_Lagrange = Lagrangian_Method(A, b, C, d);

% Display the Frobenius norm of the difference between the two solutions
disp("Gauss 消去法得到的解和精确解 (Lagrange 乘子法得到的解) 之差的 Frobenius 范数:");
disp(norm(x_Gauss - x_Lagrange, "fro"));

```

运行结果:

```

Gauss 消去法得到的解和精确解 (Lagrange 乘子法得到的解) 之差的 Frobenius 范数:
2.5071e-10

```

这验证了上述算法的正确性.

Problem 03

The least squares problem $\min \|Ax - b\|_2^2$ is equivalent to an augmented linear system:

$$\begin{bmatrix} I_{m \times m} & A \\ A^H & 0_{n \times n} \end{bmatrix} \begin{bmatrix} r \\ x \end{bmatrix} = \begin{bmatrix} b \\ 0_n \end{bmatrix}$$

which is Hermitian and indefinite.

Similar augmented linear systems exist for the linearly constrained least squares problem

$$\min \{\|Ax - b\|_2^2 : Cx = d\}$$

Can you figure it out?

- **TA:** 第三题增广系统可以进一步改写到不含有 $A^H A$: (也是由 KKT 条件导出的)

$$\begin{cases} \nabla_x L(x, \lambda) = -2A^H b + 2A^H A x + C^H \lambda = 0_n \\ Cx = d \end{cases} \Rightarrow \begin{cases} r + Ax = b \\ A^H r + C^H \lambda = 0_n \\ Cx = d \end{cases} \Rightarrow \begin{bmatrix} I_{m \times m} & A & \\ A^H & 0_{n \times n} & C^H \\ C & & 0_{p \times p} \end{bmatrix} \begin{bmatrix} r \\ x \\ \lambda \end{bmatrix} = \begin{bmatrix} b \\ 0_n \\ d \end{bmatrix}$$

Solution:

考虑线性等式约束的最小二乘问题:

$$\min_{Cx=d} \|Ax - b\|_2^2 \text{ where } \begin{cases} A \in \mathbb{C}^{m \times n} \text{ and } \text{rank}(A) = n \leq m \\ b \in \mathbb{C}^m \\ C \in \mathbb{C}^{p \times n} \text{ and } \text{rank}(C) = p < n \\ d \in \mathbb{C}^p \end{cases}$$

注意到目标函数 $f(x) = \|Ax - b\|_2^2$ 是关于 x 的凸函数, 而问题只有线性等式约束 $Cx = d$ 因此这是一个标准形式的凸优化问题, 其最优解即为 KKT 点.

定义其 Lagrange 函数 $L(x, \lambda)$ 为:

$$\begin{aligned} L(x, \lambda) &= f(x) + \lambda^H(Cx - d) \\ &= \|Ax - b\|_2^2 + \lambda^H(Cx - d) \\ \text{dom}\{L\} &= \mathbb{C}^n \times \mathbb{C}^p \end{aligned}$$

Lagrange 函数 $L(x, \lambda)$ 关于 x 的梯度为:

$$\begin{aligned} \nabla_x L(x, \lambda) &= \nabla_x \{\|Ax - b\|_2^2 + \lambda^H(Cx - d)\} \\ &= -A^H \cdot 2(b - Ax) + (\lambda^H C)^H \\ &= -2A^H b + 2A^H A x + C^H \lambda \end{aligned}$$

KKT 条件为:

$$\begin{cases} \nabla_x L(x, \lambda) = -2A^H b + 2A^H A x + C^H \lambda = 0_n & \textcircled{1} \\ Cx = d & \textcircled{2} \end{cases}$$

我们便可得到增广线性系统:

$$\begin{bmatrix} 2A^H A & C^H \\ C & 0_{p \times p} \end{bmatrix} \begin{bmatrix} x \\ \lambda \end{bmatrix} = \begin{bmatrix} 2A^H b \\ d \end{bmatrix}$$

更进一步, 我们可以求解上述方程组:

① 式左乘 $(A^H A)^{-1}$ 可得 $-2(A^H A)^{-1} A^H b + 2x + (A^H A)^{-1} C^H \lambda = 0_n$

于是有 $x = (A^H A)^{-1} A^H b - \frac{1}{2}(A^H A)^{-1} C^H \lambda$

代入 ② 式即得 $Cx = C(A^H A)^{-1} A^H b - \frac{1}{2}C(A^H A)^{-1} C^H \lambda = d$

解得 $\lambda_{\text{KKT}} = 2[C(A^H A)^{-1} C^H]^{-1}[C(A^H A)^{-1} A^H b - d]$

因此我们有:

$$\begin{aligned} x_{\text{Lagrange}} &= x_{\text{KKT}} \\ &= (A^H A)^{-1} A^H b - \frac{1}{2}(A^H A)^{-1} C^H \lambda_{\text{KKT}} \\ &= (A^H A)^{-1} A^H b - \frac{1}{2}(A^H A)^{-1} C^H \cdot 2[C(A^H A)^{-1} C^H]^{-1}[C(A^H A)^{-1} A^H b - d] \\ &= x_{\text{ls}} - (A^H A)^{-1} C^H [C(A^H A)^{-1} C^H]^{-1}(Cx_{\text{ls}} - d) \end{aligned}$$

其中 $x_{\text{ls}} := (A^H A)^{-1} A^H b$

即上述增广线性系统的解为:

$$\begin{bmatrix} x_{\text{KKT}} \\ \lambda_{\text{KKT}} \end{bmatrix} = \begin{bmatrix} x_{\text{ls}} - (A^H A)^{-1} C^H [C(A^H A)^{-1} C^H]^{-1}(Cx_{\text{ls}} - d) \\ 2[C(A^H A)^{-1} C^H]^{-1}[C(A^H A)^{-1} A^H b - d] \end{bmatrix}$$

Problem 04

Implement the Arnoldi process based on CGS/CGS2/MGS/MGS2 and visualize the loss of orthogonality.

You can randomly generate a 1000×1000 matrix and generate a 30-dimensional Krylov subspace.

(Optional) Design an algorithm to perform the Arnoldi process based on Householder orthogonalization.

Hint: use the left-looking variant.

Algorithm 1

记 $Q := [q_1, \dots, q_r]$ (其中 $1 \leq r \leq \text{rank}(Y)$)

根据 $AQ = QH$ (其中 $H \in \mathbb{C}^{r \times r}$ 为 Hessenberg 阵), 因而有 $Aq_j = \sum_{i=1}^{j+1} h_{i,j} q_i$

由于 Q 列标准正交, 故我们有 $q_k^H A q_j = \sum_{i=1}^{j+1} h_{i,j} q_k^H q_i = 0 + h_{k,j} \cdot 1 = h_{k,j}$ ($1 \leq k \leq j$)

最后有 $h_{j+1,j} q_{j+1} = Aq_j - \sum_{i=1}^j h_{i,j} q_i$

得到 $AQ = QH + \text{rank-one}$

其中 rank-one 代表遗留的秩一矩阵, 仅在最后一列有非零元素.

它的存在是因为 Aq_r 不一定能完全表示为 q_1, \dots, q_r 的线性组合.

$$A[q_1, q_2, \dots, q_r] = [Aq_1, Aq_2, \dots, Aq_r] = [q_1, q_2, \dots, q_r] \begin{bmatrix} h_{11} & h_{21} & \cdots & h_{1r} \\ h_{21} & h_{22} & \cdots & h_{2r} \\ \ddots & \ddots & \ddots & \vdots \\ & & h_{r,r-1} & h_{r,r} \end{bmatrix} + \text{rank-one}$$

基于 Gram-Schmidt 正交化的 Arnoldi 过程:

Given square matrix $A \in \mathbb{C}^{n \times n}$, vector $b \in \mathbb{C}^n$ and integer $1 \leq r \leq n$

function: $[Q, H] = \text{Gram_Schmidt_Arnoldi}(A, b, \text{tolerance}, \text{modified}, \text{reorthogonalized})$

```

n = dim(A)
Q = zeros(n, n)
Q(1 : n, 1) = b / ||b||_2
H = zeros(n, n)
delta = zeros(n, 1)

```

```

if reorthogonalized == TRUE
    max_iter = 2
else
    max_iter = 1
end

```

```

for k = 1 : r - 1
    Q(1 : n, k + 1) = A * Q(1 : n, k)
    if modified == TRUE (MGS: Modified Gram-schmidt)
        for iter = 1:max_iter
            for i = 1 : k
                delta(i) = Q(1 : n, i)^H * Q(1 : n, k + 1)
                H(i, k) = H(i, k) + delta(i)
                Q(1 : n, k + 1) = Q(1 : n, k + 1) - delta(i) * Q(1 : n, i)
            end
        end
    else (CGS: Classic Gram-Schmidt)
        for iter = 1:max_iter
            delta(1 : k) = Q(1 : n, 1 : k)^H * Q(1 : n, k + 1)
            H(1 : k, k) = H(1 : k, k) + delta(1 : k)
            Q(1 : n, k + 1) = Q(1 : n, k + 1) - Q(1 : n, 1 : k) * delta(1 : k)
        end
    end
    H(k + 1, k) = ||Q(1 : n, k + 1)||_2
    if H(k + 1, k) < tolerance (indicates linear dependence)
        r = k
        break
    else
        Q(1 : n, k + 1) = 1 / H(k + 1, k) * Q(1 : n, k + 1)
    end

```

```

    end

```

```

H(1 : r, r) = Q(1 : n, 1 : r)^H * (A * Q(1 : n, r)) (fill the last column)

```

```

Q = Q(1 : n, 1 : r)
H = H(1 : r, 1 : r)
end

```

其 Matlab 代码为:

```
function [Q, H] = Gram_Schmidt_Arnoldi(A, b, r, tolerance, modified, reorthogonalized)
```

```

% Gram_Schmidt_Arnoldi computes an orthonormal basis Q and a Hessenberg matrix H
% using the Arnoldi process with either Classical or Modified Gram-Schmidt.
%
% Inputs:
%   A - Square matrix of size n x n.
%   b - Initial vector of size n x 1.
%   r - Desired rank of the output.
%   tolerance - Threshold for detecting linear dependence (default: 1e-10).
%   modified - Flag for using Modified Gram-Schmidt (default: false).
%   reorthogonalized - Flag for reorthogonalization (default: false).
%
% Outputs:
%   Q - Orthonormal basis of size n x r.
%   H - Upper Hessenberg matrix of size r x r.

% Validate inputs and set default values if not provided
if nargin < 6
    reorthogonalized = false; % Default to not reorthogonalizing
end
if nargin < 5
    modified = false; % Default to Classical Gram-Schmidt
end
if nargin < 4
    tolerance = 1e-10; % Default tolerance for linear dependence
end

% Get the size of the matrix A
n = size(A, 1);

% Check if the desired rank is valid
if r < 1 || r > n
    error("r should be an integer in [1, n]");
end

% Initialize matrices Q and H
Q = zeros(n, n);
Q(:, 1) = b / norm(b); % Normalize the initial vector b
H = zeros(n, n); % Initialize H to zeros
delta = zeros(n, 1); % Temporary variable for inner products

% Set max iterations based on reorthogonalization flag
if reorthogonalized
    max_iter = 2; % More iterations for reorthogonalization
else
    max_iter = 1; % Standard iteration count
end

% Loop to build the orthonormal basis up to r-1 or n-1
for k = 1:r-1
    % Apply the matrix A to the last basis vector
    Q(:, k + 1) = A * Q(:, k);

    if modified
        % Modified Gram-Schmidt process
        for iter = 1:max_iter
            for i = 1:k
                % Compute inner product
                delta(i) = Q(:, i)' * Q(:, k + 1);
                % Update H matrix
                H(i, k) = H(i, k) + delta(i);
            end
            % Orthogonalize the k+1 vector
            for i = 1:k
                Q(:, k + 1) = Q(:, k + 1) - delta(i) * Q(:, i);
            end
            delta(k + 1) = norm(Q(:, k + 1));
            if delta(k + 1) < tolerance
                break;
            end
        end
    end
end

```

```

        Q(:, k + 1) = Q(:, k + 1) - delta(i) * Q(:, i);
    end
end
else
    % Classical Gram-Schmidt process
    for iter = 1:max_iter
        % Compute inner products for all previous basis vectors
        delta(1:k) = Q(:, 1:k)' * Q(:, k + 1);
        % Update H matrix
        H(1:k, k) = H(1:k, k) + delta(1:k);
        % Orthogonalize the k+1 vector
        Q(:, k + 1) = Q(:, k + 1) - Q(:, 1:k) * delta(1:k);
    end
end

% Compute the norm for the current basis vector
H(k + 1, k) = norm(Q(:, k + 1));

% Check for linear dependence by comparing the norm with the tolerance
if H(k + 1, k) < tolerance
    fprintf("The rank %d is lesser than %d\n", k, r);
    r = k; % Update the rank if linear dependence is detected
    break; % Exit the loop early
else
    % Normalize the current basis vector
    Q(:, k + 1) = Q(:, k + 1) / H(k + 1, k);
end
end

% If the rank is full, fill the last column of H
H(1:r, r) = Q(:, 1:r)' * (A * Q(:, r));

% Trim Q and H to the computed effective rank
Q = Q(:, 1:r);
H = H(1:r, 1:r);
end

```

可视化正交性损失的函数 `visualize_orthogonality_loss`:

```

function visualize_orthogonality_loss(Q, titleStr)
    % Visualizes the componentwise loss of orthogonality |Q^H Q - I_n|
    loss = Q' * Q - eye(size(Q, 2)); % Compute the loss
    figure; % Create a new figure window
    imagesc(log10(abs(loss))); % Display the absolute value of the loss
    colorbar; % Add colorbar to indicate scale
    title(titleStr);
    xlabel('Column Index');
    ylabel('Row Index');
    axis square; % Make the axes square for better visualization
end

```

函数调用:

```

rng(51);
n = 1000; % Size of the square matrix
r = 30; % Dimension of the Krylov subspace

% Generate a random complex matrix A and vector b
A = rand(n, n) + 1i * rand(n, n);
b = rand(n, 1) + 1i * rand(n, 1);

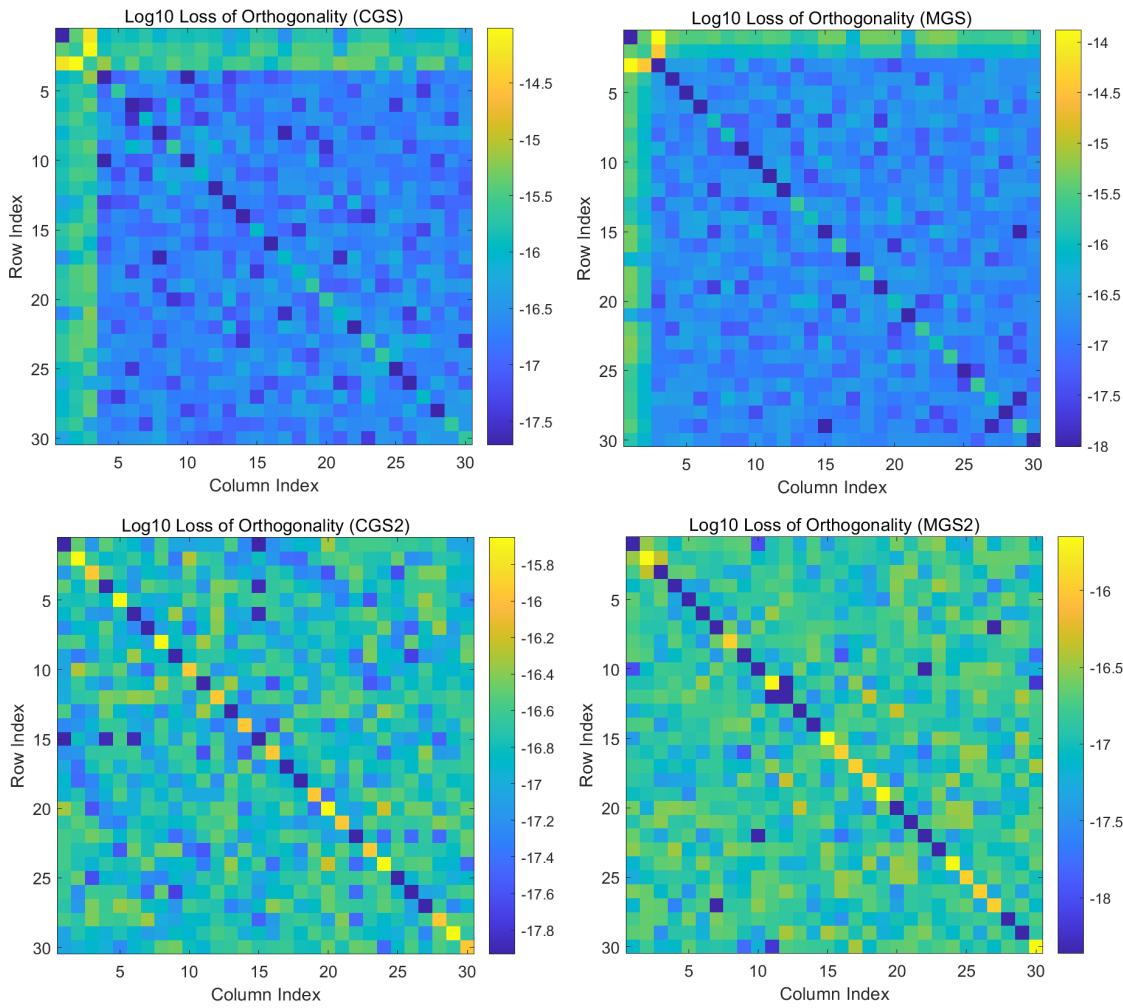
```

```
% Apply the Arnoldi process with different methods
[Q_CGS, H_CGS] = Gram_Schmidt_Arnoldi(A, b, r, 1e-10, false, false);
visualize_orthogonality_loss(Q_CGS, 'Log10 Loss of Orthogonality (CGS)');
disp("CGS:")
leftover_CGS = A * Q_CGS(:, end) - Q_CGS(:, 1:end) * H_CGS(1:end, end);
residual_CGS = A * Q_CGS - Q_CGS * H_CGS;
residual_CGS(:, end) = residual_CGS(:, end) - leftover_CGS;
disp(norm(residual_CGS, "fro") / norm(A, "fro"));

[Q_CGS2, H_CGS2] = Gram_Schmidt_Arnoldi(A, b, r, 1e-10, false, true);
visualize_orthogonality_loss(Q_CGS2, 'Log10 Loss of Orthogonality (CGS2)');
disp("CGS2:")
leftover_CGS2 = A * Q_CGS2(:, end) - Q_CGS2(:, 1:end) * H_CGS2(1:end, end);
residual_CGS2 = A * Q_CGS2 - Q_CGS2 * H_CGS2;
residual_CGS2(:, end) = residual_CGS2(:, end) - leftover_CGS2;
disp(norm(residual_CGS2, "fro") / norm(A, "fro"));

[Q_MGS, H_MGS] = Gram_Schmidt_Arnoldi(A, b, r, 1e-10, true, false);
visualize_orthogonality_loss(Q_MGS, 'Log10 Loss of Orthogonality (MGS)');
disp("MGS:")
leftover_MGS = A * Q_MGS(:, end) - Q_MGS(:, 1:end) * H_MGS(1:end, end);
residual_MGS = A * Q_MGS - Q_MGS * H_MGS;
residual_MGS(:, end) = residual_MGS(:, end) - leftover_MGS;
disp(norm(residual_MGS, "fro") / norm(A, "fro"));

[Q_MGS2, H_MGS2] = Gram_Schmidt_Arnoldi(A, b, r, 1e-10, true, true);
visualize_orthogonality_loss(Q_MGS2, 'Log10 Loss of Orthogonality (MGS2)');
disp("MGS2:")
leftover_MGS2 = A * Q_MGS2(:, end) - Q_MGS2(:, 1:end) * H_MGS2(1:end, end);
residual_MGS2 = A * Q_MGS2 - Q_MGS2 * H_MGS2;
residual_MGS2(:, end) = residual_MGS2(:, end) - leftover_MGS2;
disp(norm(residual_MGS2, "fro") / norm(A, "fro"));
```



Algorithm 2

记 $Q := [q_1, \dots, q_r]$ (其中 $1 \leq r \leq \text{rank}(Y)$)

Arnoldi 过程的基本框架是 $AQ = QH + \text{rank-one}$

其中 $H \in \mathbb{C}^{r \times r}$ 为 Hessenberg 阵，而 rank-one 代表遗留的秩一矩阵，仅在最后一列有非零元素.

它的存在是因为 Aq_r 不一定能完全表示为 q_1, \dots, q_r 的线性组合.

我们可以通过扩充一列，将 H 扩充为上三角阵，使用 Householder QR 做:

$$A[q_1, q_2, \dots, q_r] = [Aq_1, Aq_2, \dots, Aq_r] = [q_1, q_2, \dots, q_r] \begin{bmatrix} h_{11} & h_{21} & \cdots & h_{1r} \\ h_{21} & h_{22} & \cdots & h_{2r} \\ \ddots & \ddots & \ddots & \vdots \\ & & h_{r,r-1} & h_{r,r} \end{bmatrix} + \text{rank-one}$$

$$[q_1, Aq_1, Aq_2, \dots, Aq_r] = [q_1, q_2 \dots, q_r] \begin{bmatrix} 1 & h_{11} & h_{21} & \cdots & h_{1r} \\ h_{21} & h_{22} & \cdots & h_{2r} \\ \ddots & \ddots & \ddots & \vdots \\ & & h_{r,r-1} & h_{r,r} \end{bmatrix} + \text{rank-one}$$

值得注意的是，我们并不是直接生成 $[q_1, Aq_1, Aq_2, \dots, Aq_r]$ 然后应用 Householder QR

而是使用 **left-looking** 的变体，即存储每步的 Householder 变换

(具体来说可使用 **Matrix Computation 算法 5.1.2**,

但注意其 Householder 变换的计算顺序与实际作用在新加入的向量的顺序相反)

当添加新的向量时，对其应用之前的 Householder 变换，再计算该步需要进行的 Householder 变换.

(Matrix Computation 算法 5.1.2)

设有 $r \leq n$ 个 Householder 变换 H_1, \dots, H_r ，其中:

$$H_j = I_n - \beta_j v^{(j)} (v^{(j)})^H$$

$$v^{(j)} = [\underbrace{0, \dots, 0}_{j-1}, 1, v_{j+1}^{(j)}, \dots, v_n^{(j)}]^T$$

我们可以计算得到 $W, Y \in \mathbb{C}^{n \times r}$ 满足 $H_1 \cdots H_r = I_n + WY^H$:

$$\begin{aligned} Y &= v^{(1)} \\ W &= -\beta_1 v^{(1)} \\ \text{for } j &= 2 : r \\ z &= -\beta_j (I_n + WY^H) v^{(j)} = -\beta_j [v^{(j)} + W(Y^H v^{(j)})] \\ W &= [W, z] \\ Y &= [Y, v^{(j)}] \\ \text{end} \end{aligned}$$

复数域上使用 Householder 变换计算 Arnoldi 过程的函数:

(注意其中应用 $(I - WY^H)^H = I - YW^H$ 左乘新加入的向量, 而不是用 $I - WY$ 左乘新加入的向量, 这个 bug 找了很久)

```
function [Q, H] = Complex_Householder_Arnoldi(A, b, r, tolerance)
    % This function performs the Complex Householder Arnoldi process.
    % Input:
    %   A           : The input matrix (n x n)
    %   b           : The input vector (n x 1)
    %   r           : The number of desired orthonormal vectors
    %   tolerance : Tolerance for stopping criterion
    % Output:
    %   Q           : Orthonormal basis of the Krylov subspace (n x r)
    %   H           : Upper Hessenberg matrix (r x (r+1))

    n = size(A, 1); % Get the size of matrix A (assumes A is square)

    % Initialize matrices
    Y = zeros(n, r); % Matrix to store Householder vectors
    W = zeros(n, r); % Matrix to store Householder weights
    H = zeros(n, r + 1); % Matrix to store the upper Hessenberg matrix
    Q = zeros(n, r); % Matrix to store the orthonormal basis vectors

    % Main loop to construct the Arnoldi process
    for k = 1:r
        if k == 1
            % First iteration: Apply Householder transformation to b
            [v, beta] = Complex_Householder(b); % Get the Householder vector and scalar
            beta
            Y(:, 1) = v; % Store the Householder vector in Y
            W(:, 1) = -beta * v; % Store the modified vector W
            v = zeros(n, 1); % Reset v for future use

            H(1, 1) = norm(b, 2) * b(1) / abs(b(1)); % Store the norm of b in H(1,1)
            Q(:, 1) = W(:, 1) * Y(1, 1)'; % Compute the first orthonormal vector
        for Q
            Q(1, 1) = Q(1, 1) + 1; % Adjust the first element of Q to
            maintain orthonormality

        else
            % Subsequent iterations
            z = A * Q(:, k-1); % Matrix-vector product for the next
            Krylov subspace vector
```

```

z = z + Y(:, 1:k-1) * (W(:, 1:k-1)' * z); % Adjust z with contributions
from W and Y

% Store the computed values in H
H(1:k-1, k) = z(1:k-1); % Fill the upper part of the current
column in H
H(k, k) = norm(z(k:n), 2) * z(k) / abs(z(k)); % Compute the norm of the
remainder of z and store in H

% Check for convergence based on the tolerance
if abs(H(k, k)) < tolerance
    r = k - 1; % Adjust the rank if the value is below
tolerance
    break % Exit the loop if convergence is
achieved
end

% Apply Householder transformation to the part of z
[v(k:n), beta] = Complex_Householder(z(k:n)); % Generate the Householder
vector and beta for z(k:n)
Y(:, k) = v; % Store the Householder vector in Y
W(:, k) = -beta * (v + W(:, 1:k-1) * (Y(:, 1:k-1)' * v)); % Update W based
on the new Householder vector
v = zeros(n, 1); % Reset v after use
Q(:, k) = W(:, 1:k) * Y(k, 1:k)'; % Compute the k-th orthonormal vector
for Q
    Q(k, k) = Q(k, k) + 1; % Adjust the diagonal of Q to
maintain orthonormality
end
end

% If the rank r is full, compute the last column of H
z = A * Q(:, r); % Matrix-vector product for the last vector in Q
z = z + W(:, 1:r) * (Y(:, 1:r)' * z); % Adjust z with contributions from W and Y
H(1:r, r+1) = Q(:, 1:r)' * z; % Store the last column in H by projecting z onto Q

% Trim H and Q to their final sizes based on the computed rank
H = H(1:r, 2:r + 1); % Keep only the relevant portion of H
Q = Q(:, 1:r); % Keep only the relevant portion of Q
end

```

复数域上的 Householder 变换的计算算法已在 Homework 4 Problem 2 中给出:

```

function [v, beta] = Complex_Householder(x)
% This function computes the Householder vector 'v' and scalar 'beta' for
% a given complex vector 'x'. This transformation is used to create zeros
% below the first element of 'x' by reflecting 'x' along a specific direction.

n = length(x);
x = x / norm(x, inf); % Normalize x by its infinity norm to avoid numerical issues

% Copy all elements of 'x' except the first into 'v'
v = zeros(n, 1);
v(2:n) = x(2:n);

% Compute sigma as the squared 2-norm of the elements of x starting from the second
element
sigma = norm(x(2:n), 2)^2;

% Check if sigma is near zero, which would mean 'x' is already close to a scalar
multiple of e_1

```

```

if sigma < 1e-10
    beta = 0; % If sigma is close to zero, set beta to zero (no transformation
needed)
else
    % Determine gamma to account for the argument of complex number x(1)
    if abs(x(1)) < 1e-10
        gamma = 1; % If x(1) is close to zero, set gamma to 1
    else
        gamma = x(1) / abs(x(1)); % otherwise, set gamma to x(1) divided by its
magnitude
    end

    % Compute alpha as the Euclidean norm of x, including x(1) and sigma
    alpha = sqrt(abs(x(1))^2 + sigma);

    % Compute the first element of 'v' to avoid numerical cancellation
    v(1) = -gamma * sigma / (abs(x(1)) + alpha);

    % Calculate 'beta', the scaling factor of the Householder transformation
    beta = 2 * abs(v(1))^2 / (abs(v(1))^2 + sigma);

    % Normalize the vector 'v' by v(1) to ensure that the first element is 1,
    % allowing for simplified storage and computation of the transformation
    v = v / v(1);
end
end

```

函数调用:

```

rng(51);
n = 1000; % Size of the square matrix
r = 30; % Dimension of the Krylov subspace
tolerance = 1e-10;

% Generate a random complex matrix A (n x n) and vector b (n x 1)
A = rand(n, n) + 1i * rand(n, n); % Create a random complex matrix with real and
imaginary parts
b = rand(n, 1) + 1i * rand(n, 1); % Create a random complex vector with real and
imaginary parts

% Call the Complex Householder Arnoldi process to compute Q and H
[Q, H] = Complex_Householder_Arnoldi(A, b, r, tolerance);

% visualize the loss of orthogonality in the Q matrix
visualize_orthogonality_loss(Q, 'Log10 Loss of Orthogonality (Householder-Arnoldi)');

% Compute the "leftover" term for the last column of the residual
leftover = A * Q(:, end) - Q(:, 1:end) * H(1:end, end);

% Calculate the residual matrix between AQ and QH, ensuring the last column accounts
for the leftover term
residual = A * Q - Q * H;
residual(:, end) = residual(:, end) - leftover; % Adjust the last column with the
leftover term

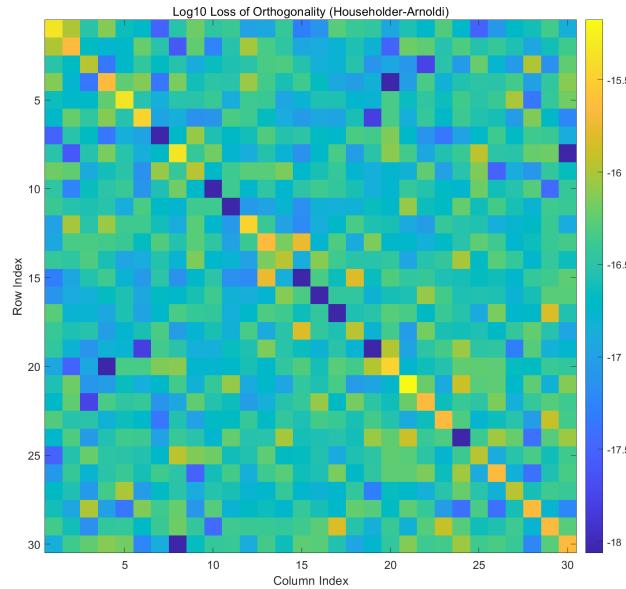
% Display the Frobenius norm of the residual matrix, normalized by the Frobenius norm
of A
disp("Householder (left-looking):")
disp(norm(residual, "fro") / norm(A, "fro"));

```

运行结果:

Householder (left-looking):

9.4769e-16



Problem 05

Write a program to solve rank deficient least squares problems.

You may assume that the rank is given.

Test your program with a low-rank least squares problem.

Try it with a general full-rank least squares solver and compare the results.

- TA: 第五题也可以用选主元的 QR 分解求解秩亏损的最小二乘问题

Solution:

(1) Helper Functions

Gram-Schmidt QR 分解算法已经在 Homework 5 Problem 1 中给出了:

```
function [Q, R] = Gram_Schmidt_QR(A, tolerance, modified, reorthogonalized)
    % This function performs the Gram-Schmidt QR factorization of a matrix A
    % It supports both classical and modified versions of the GS algorithm,
    % and it allows for reorthogonalization to improve numerical stability.
    %
    % Inputs:
    %   - A: The m x n matrix to be factorized
    %   - tolerance: The threshold below which a vector is considered linearly
    %     dependent
    %   - modified: Boolean flag to choose between Classical Gram-Schmidt (CGS)
    %               or Modified Gram-Schmidt (MGS)
    %   - reorthogonalized: Boolean flag to perform reorthogonalization (improves
    %     numerical stability)
    %
    % Outputs:
    %   - Q: An m x r orthonormal matrix (r is the rank of A, or the number of
    %     orthogonal vectors)
    %   - R: An r x n upper triangular matrix
```

```

[m, n] = size(A); % Get the size of matrix A (m rows, n columns)
r = 0; % Initialize rank of A
Q = zeros(m, m); % Preallocate Q as an m x m zero matrix
R = zeros(m, n); % Preallocate R as an m x n zero matrix
delta = zeros(m, 1); % Temporary vector for storing projection coefficients

% Set the number of orthogonalization iterations based on the reorthogonalized flag
if reorthogonalized
    max_iter = 2; % If reorthogonalization is enabled, perform two passes
else
    max_iter = 1; % Otherwise, perform only one pass
end

% Main loop over each column of matrix A (for each column k)
for k = 1:min(m,n)
    % Initialize the k-th column of Q as the k-th column of A
    Q(1:m, r+1) = A(1:m, k);

    % If modified Gram-Schmidt (MGS) is selected
    if modified
        for iter = 1:max_iter % Repeat orthogonalization based on max_iter
            for i = 1:r % Loop over previously computed columns of Q
                delta(i) = Q(1:m, i)' * Q(1:m, r+1); % Compute projection of Q_k
on Q_i
                R(i, k) = R(i, k) + delta(i); % Update the corresponding entry in
R
                Q(1:m, r+1) = Q(1:m, r+1) - delta(i) * Q(1:m, i); % Subtract
projection from Q_k
            end
        end
    else % Classical Gram-Schmidt (CGS)
        for iter = 1:max_iter
            delta(1:r) = Q(1:m, 1:r)' * Q(1:m,r+1); % Compute projections in one
step
            R(1:r, k) = R(1:r, k) + delta(1:r); % Update R
            Q(1:m, r+1) = Q(1:m, r+1) - Q(1:m, 1:r) * delta(1:r); % Subtract the
projection from Q_k
        end
    end

    % Compute the 2-norm of the current column of Q (for normalization)
    R(r+1, k) = norm(Q(1:m, r+1), 2);

    % Check if the norm is smaller than the tolerance, indicating linear dependence
    if R(r+1, k) < tolerance
        R(r+1, k) = 0; % Set R entry to zero if linearly dependent
    else
        Q(1:m, r+1) = Q(1:m, r+1) / R(r+1, k); % Normalize the vector
        r = r + 1; % Increment the rank
    end

end

% Additional step: if the number of columns n is greater than m,
% compute the remaining upper triangular part of R using the orthonormal Q matrix
if n > m
    for k = m+1:n
        % Compute the projections of columns of A onto the previously computed
orthonormal columns of Q
        R(1:r, k) = Q(1:m, 1:r)' * A(1:m, k);
    end
end

```

```

        end
    end

    % Reduce the size of Q and R to the actual rank r of A
    Q = Q(1:m, 1:r); % Return the first r columns of Q
    R = R(1:r, 1:n); % Return the first r rows of R
end

```

`Rotate()` 函数用于将一个矩阵绕矩阵中心旋转 180° ，最后再进行简单转置(即不取复共轭):

```

function B = Rotate(A)
    % Rotate computes the 90-degree counterclockwise rotation of a matrix A.
    %
    % Inputs:
    %   A - The input matrix to be rotated (size m x n).
    %
    % Outputs:
    %   B - The rotated matrix (size n x m).

    [m, n] = size(A); % Get the dimensions of the input matrix A

    B = zeros(n, m); % Initialize the output matrix B with zeros (size n x m)

    % Loop through each element of A to fill in the rotated matrix B
    for i = 1:m
        for j = 1:n
            % Place the element from A into its new position in B
            B(n - j + 1, m - i + 1) = A(i, j);
        end
    end

    % Note: In MATLAB, the operation A' is the complex conjugate transpose,
    % while A.' is the simple transpose.
    % In this function, B = B.';
    % ensures we only transpose the matrix without taking the complex conjugate.
    B = B.';

    % Transpose the resulting matrix to correct the orientation
end

```

回代法:

```

function X = Backward_Sweep(U, Y)
    % 回代法求解 UX = Y
    [n, ~] = size(Y);
    for i = n:-1:2
        Y(i,:) = Y(i,:) / U(i, i); % 对角线归一化
        Y(1:i-1,:) = Y(1:i-1,:) - U(1:i-1, i) * Y(i,:);
    end
    Y(1,:) = Y(1,:) / U(1, 1); % 处理第一行
    X = Y; % 返回结果
end

```

(2) 基于 QR 的 RQ 算法

基于已有的 QR 算法，我们可以方便地得到 QL 分解，进而得到 RQ 分解:

上述算法的 Matlab 代码为:

```

function [R, Q] = Gram_Schmidt_RQ(A, tolerance, modified, reorthogonalized)
    % Gram_Schmidt_RQ computes the RQ factorization of matrix A using
    % the Gram-Schmidt process.
    %
    % Inputs:
    %   A - The input matrix to be factorized (size m x n).
    %   tolerance - Threshold for the modified Gram-Schmidt process.
    %   modified - Boolean indicating whether to use modified Gram-Schmidt.
    %   reorthogonalized - Boolean indicating whether to reorthogonalize.
    %
    % Outputs:
    %   R - Upper triangular matrix (size n x n).
    %   Q - Orthogonal matrix (size m x n).

    A = A'; % Transpose A to prepare for RQ factorization

    % Call the Gram-Schmidt QR function on the reversed matrix
    [Q, R] = Gram_Schmidt_QR(A(:, end:-1:1), tolerance, modified, reorthogonalized);

    Q = Q(:, end:-1:1); % Reverse the columns of Q to match RQ format
    R = Rotate(R); % Apply rotation to R for final format

    Q = Q'; % Transpose Q back to original dimensions
    R = R'; % Transpose R back to original dimensions
end

```

(3) 亏秩最小二乘问题的求解

考虑亏秩最小二乘问题 $\min_x \|Ax - b\|_2^2$

其中 $b \in \mathbb{C}^m$ 为给定向量, 而 $A \in \mathbb{C}^{m \times n}$ 为秩 $r < \min(m, n)$ 的矩阵.

假设我们已通过一次 QR 分解和一次 RQ 分解得到了秩 r 矩阵 $A \in \mathbb{C}^{m \times n}$ 的分解 $A = Q_1 R Q_2$
 其中 $Q_1 \in \mathbb{C}^{m \times r}$ 为列标准正交的矩阵, $R \in \mathbb{C}^{r \times r}$ 为满秩的上三角阵, $Q_2 \in \mathbb{C}^{r \times n}$ 为行标准正交的矩阵.
 则我们有:

$$\begin{aligned}
\min_x \|Ax - b\|_2^2 &\Leftrightarrow \min_x \|Q_1 R Q_2 x - b\|_2^2 \\
&\Leftrightarrow \min_x \left\| \begin{bmatrix} Q_1 & Q_1^\perp \end{bmatrix} \begin{bmatrix} R \\ 0_{(m-r) \times r} \end{bmatrix} Q_2 x - b \right\|_2^2 \\
&\Leftrightarrow \min_x \left\| \begin{bmatrix} R \\ 0_{(m-r) \times r} \end{bmatrix} Q_2 x - \begin{bmatrix} Q_1^H b \\ (Q_1^\perp)^H b \end{bmatrix} \right\|_2^2 \\
&\Leftrightarrow \min_x \|R Q_2 x - Q_1^H b\|_2^2 \\
&\Leftrightarrow \min_x \|Q_2 x - R^{-1} Q_1^H b\|_2^2
\end{aligned}$$

因此 $x = Q_2^H R^{-1} Q_1^H b$ 是亏秩最小二乘问题 $\min_x \|Ax - b\|_2^2$ 的一个特解.

下面我们给出求解亏秩最小二乘问题 $\min_x \|Ax - b\|_2^2$ 的特解 $x = Q_2^H R^{-1} Q_1^H b$ 的 Matlab 函数:

```

function [x, r] = Rank_Deficient_Solver(A, b)
    % Solve_Rank_Deficient_Least_Squares computes the least squares solution
    % for a potentially rank-deficient matrix A using the Gram-Schmidt process.
    %
    % Inputs:
    %   A - The input matrix (size m x n), which may be rank deficient.
    %   b - The right-hand side vector (size m x 1).
    %
    % Outputs:
    %   x - The least squares solution vector (size n x 1).
    %   r - The rank of the matrix R obtained during the process.

    % Step 1: Perform QR factorization of A using the Gram-Schmidt process
    [Q1, R_tmp] = Gram_Schmidt_QR(A, 1e-10, true, true);

    % Step 2: Get the size of R_tmp to determine the effective rank
    [r, ~] = size(R_tmp); % r will be the number of rows in R_tmp

    % Step 3: Perform RQ factorization on R_tmp
    [R, Q2] = Gram_Schmidt_RQ(R_tmp, 1e-10, true, true);

    % Step 4: Solve the triangular system R * b_tilde = Q1' * b
    b_tilde = Backward_Sweep(R, Q1' * b); % Backward substitution

    % Step 5: Compute the final solution x using Q2
    x = Q2' * b_tilde; % Project the solution back

end

```

(4) 满秩求解器

复数域上的 Householder 变换的计算算法已在 Homework 4 Problem 2 中给出:

```

function [v, beta] = Complex_Householder(x)
    % This function computes the Householder vector 'v' and scalar 'beta' for
    % a given complex vector 'x'. This transformation is used to create zeros
    % below the first element of 'x' by reflecting 'x' along a specific direction.

    n = length(x);
    x = x / norm(x, inf); % Normalize x by its infinity norm to avoid numerical issues

    % Copy all elements of 'x' except the first into 'v'
    v = zeros(n, 1);
    v(2:n) = x(2:n);

```

```

% Compute sigma as the squared 2-norm of the elements of x starting from the second
element
sigma = norm(x(2:n), 2)^2;

% Check if sigma is near zero, which would mean 'x' is already close to a scalar
multiple of e_1
if sigma < 1e-10
    beta = 0; % If sigma is close to zero, set beta to zero (no transformation
needed)
else
    % Determine gamma to account for the argument of complex number x(1)
    if abs(x(1)) < 1e-10
        gamma = 1; % If x(1) is close to zero, set gamma to 1
    else
        gamma = x(1) / abs(x(1)); % otherwise, set gamma to x(1) divided by its
magnitude
    end

    % Compute alpha as the Euclidean norm of x, including x(1) and sigma
alpha = sqrt(abs(x(1))^2 + sigma);

    % Compute the first element of 'v' to avoid numerical cancellation
v(1) = -gamma * sigma / (abs(x(1)) + alpha);

    % Calculate 'beta', the scaling factor of the Householder transformation
beta = 2 * abs(v(1))^2 / (abs(v(1))^2 + sigma);

    % Normalize the vector 'v' by v(1) to ensure that the first element is 1,
    % allowing for simplified storage and computation of the transformation
v = v / v(1);
end
end

```

复数域上的 Householder QR 算法已在 Homework 4 Problem 3 中给出:

```

function [Q, R] = Complex_Householder_QR(A)
[m, n] = size(A);
Q = eye(m); % Initialize Q as the identity matrix
R = A; % Initialize R as A

for k = 1:min(m-1, n)
    [v, beta] = Complex_Householder(R(k:m, k)); % Apply Complex Householder

    % Update R
    R(k:m, k:n) = R(k:m, k:n) - (beta * v) * (v' * R(k:m, k:n));

    % Update Q
    Q(1:m, k:m) = Q(1:m, k:m) - (Q(1:m, k:m) * v) * (beta * v');
end
end

```

计算得到 $A \in \mathbb{C}^{m \times n}$ 的 QR 分解 $A = QR$ 之后 (其中 $Q \in \mathbb{C}^{m \times m}$ 为酉矩阵, $R \in \mathbb{C}^{m \times n}$ 的上 $n \times n$ 分块 R_1 为上三角阵)

求解法方程 $A^H Ax = R^H Rx = A^H b$ 就等价于求解 $\begin{cases} R^H y = A^H b \\ Rx = y \end{cases}$ (分别由前代法和回代法求解)

或者也可考虑精简 QR 分解 $A = Q_1 R_1$ (其中 $Q_1 \in \mathbb{C}^{m \times n}$ 由 Q 的前 n 列构成)

则求解法方程 $A^H Ax = R_1^H R_1 x = R_1^H Q_1^H b = A^H b$ 就等价于求解 $R_1 x = Q_1^H b$ (由回代法求解)

```

function x = Householder_Full_Rank_Solver(A, b)
[m, n] = size(A);

% Step 1: Compute the QR decomposition of A using Householder reflections
[Q, R] = Complex_Householder_QR(A);

% Step 2: Solve the system Rx = Q' * b using backward substitution
x = Backward_Sweep(R(1:n, 1:n), Q(1:m, 1:n)' * b);
end

```

(5) 运行结果

我们比较 Matlab 内置求解器、亏秩求解器和满秩求解器应用于亏秩最小二乘问题的效果.

函数调用:

```

rng(51); % Set the random seed for reproducibility

% Define dimensions
m = 300; % Number of rows in matrix A
n = 50; % Number of columns in matrix A
r = 25; % Rank of matrix A

% Generate a rank-r complex matrix A
A = (rand(m, r) + 1i * rand(m, r)) * (rand(r, n) + 1i * rand(r, n));
b = rand(m, 1) + 1i * rand(m, 1); % Generate a random complex vector b

% Ensure the rank of A is exactly r
assert(rank(A) == r);

% Solve the linear system Ax = b for the exact solution
x_exact = A \ b; % Use MATLAB's backslash operator for solving
object_value_exact = norm(A * x_exact - b, 2); % calculate the exact objective value
disp('Exact solution objective value:');
disp(object_value_exact);

% Solve for the rank-deficient case
[x_deficient, r_deficient] = Rank_Deficient_Solver(A, b); % Custom solver for rank-
deficient case
object_value_deficient = norm(A * x_deficient - b, 2); % calculate the objective value
for deficient case
disp('Rank-deficient solution objective value:');
disp(object_value_deficient);

% Solve for the full rank case using Householder transformations
x_full_rank = Householder_Full_Rank_Solver(A, b); % Custom solver for full rank case
object_value_full_rank = norm(A * x_full_rank - b, 2); % calculate the objective value
for full rank case
disp('Full rank solution objective value:');
disp(object_value_full_rank);

```

运行结果:

```

警告: 秩亏, 秩 = 25, tol = 1.703841e-10。
Exact solution objective value:
6.7824

Rank-deficient solution objective value:
6.7824

Full rank solution objective value:
7.4224

```

Problem 06 (optional)

Implement the Lanczos process.

Make sure you are using a short recurrence instead of using the naive Arnoldi process.

Solution:

当 A 是对称阵时, Arnoldi 过程得到的 $T := Q^T A Q$ 是一个对称三对角阵.

理论上, 我们可将 Arnoldi 过程简化为 **Lanczos 过程** (考虑到 T 的很多上三角元为零)

给定对称阵 $A \in \mathbb{R}^{n \times n}$ 和单位向量 $q_1 \in \mathbb{R}^n$

我们记:

$$\tilde{T}_k = \begin{bmatrix} \alpha_1 & \beta_1 & & & \\ \beta_1 & \alpha_2 & \beta_2 & & \\ & \beta_2 & \alpha_3 & \ddots & \\ & & \ddots & \ddots & \beta_{k-1} \\ & & & \beta_{k-1} & \alpha_k \\ & & & & \beta_k \end{bmatrix}$$

记 $Q_k := [q_1, \dots, q_k]$ (其中 $1 \leq k \leq \text{rank}(\mathcal{K}(A, q_1, n))$)

(其中 Krylov 子空间 $\mathcal{K}(A, q_1, n) = \text{span}\{q_1, Aq_1, \dots, A^{n-1}q_1\}$)

根据 $AQ_k = Q_{k+1}\tilde{T}_k$ 可知 $Aq_k = \beta_{k-1}q_{k-1} + \alpha_k q_k + \beta_k q_{k+1}$

由于 q_1, \dots, q_{k+1} 标准正交, 故我们有 $q_k^T A q_k = \alpha_k$

最后有 $\beta_k q_{k+1} = Aq_k - \alpha_k q_k - \beta_{k-1} q_{k-1}$

Matlab 代码如下:

```

function [Q, T] = Lanczos(A, b, r, tolerance)
    % Lanczos Algorithm to compute the tridiagonal matrix T and orthonormal basis Q
    %
    % Input:
    %   A   - A symmetric matrix (n x n)
    %   b   - A vector (n x 1)
    %   r   - Number of Lanczos iterations
    %
    % Output:
    %   Q   - Orthogonal basis vectors (n x m)
    %   T   - Tridiagonal matrix (m x m)

    % Initialize variables
    n = size(A, 1); % Get the size of matrix A
    Q = zeros(n, r); % Initialize orthogonal basis matrix Q with zeros (n x r)
    T = zeros(r, r); % Initialize tridiagonal matrix T with zeros (r x r)

    % Normalize the initial vector b to create the first orthogonal basis vector
    Q(:, 1) = b / norm(b);

```

```

% Start the Lanczos iterations
for j = 1:r
    % Compute the matrix-vector product of A and the j-th basis vector
    z = A * Q(:, j);

    % Compute the diagonal entry of T (the Rayleigh quotient)
    T(j, j) = Q(:, j)' * z;

    if j == 1
        % For the first iteration, adjust z by subtracting the first term
        z = z - T(j, j) * Q(:, j);
    else
        % For subsequent iterations, subtract the contributions from the last two
        % basis vectors
        z = z - T(j, j) * Q(:, j) - T(j-1, j) * Q(:, j-1);
    end

    % Compute the off-diagonal entry of T and check for convergence
    T(j, j+1) = norm(z, 2); % Norm of the vector z is the off-diagonal element
    T(j+1, j) = T(j, j+1); % Since T is symmetric

    if norm(z, 2) < tolerance
        % If the norm of z is below the tolerance, reduce the number of iterations
        r = j-1;
        break; % Exit the loop if convergence is achieved
    else
        % If not converged, normalize z to create the next basis vector
        Q(:, j+1) = z / T(j, j+1);
    end
end

% Return results by truncating Q and T to the size of the computed basis
Q = Q(:, 1:r); % Truncate Q to include only the computed basis vectors
T = T(1:r, 1:r); % Truncate T to the size of the computed matrix
end

```

函数调用:

```

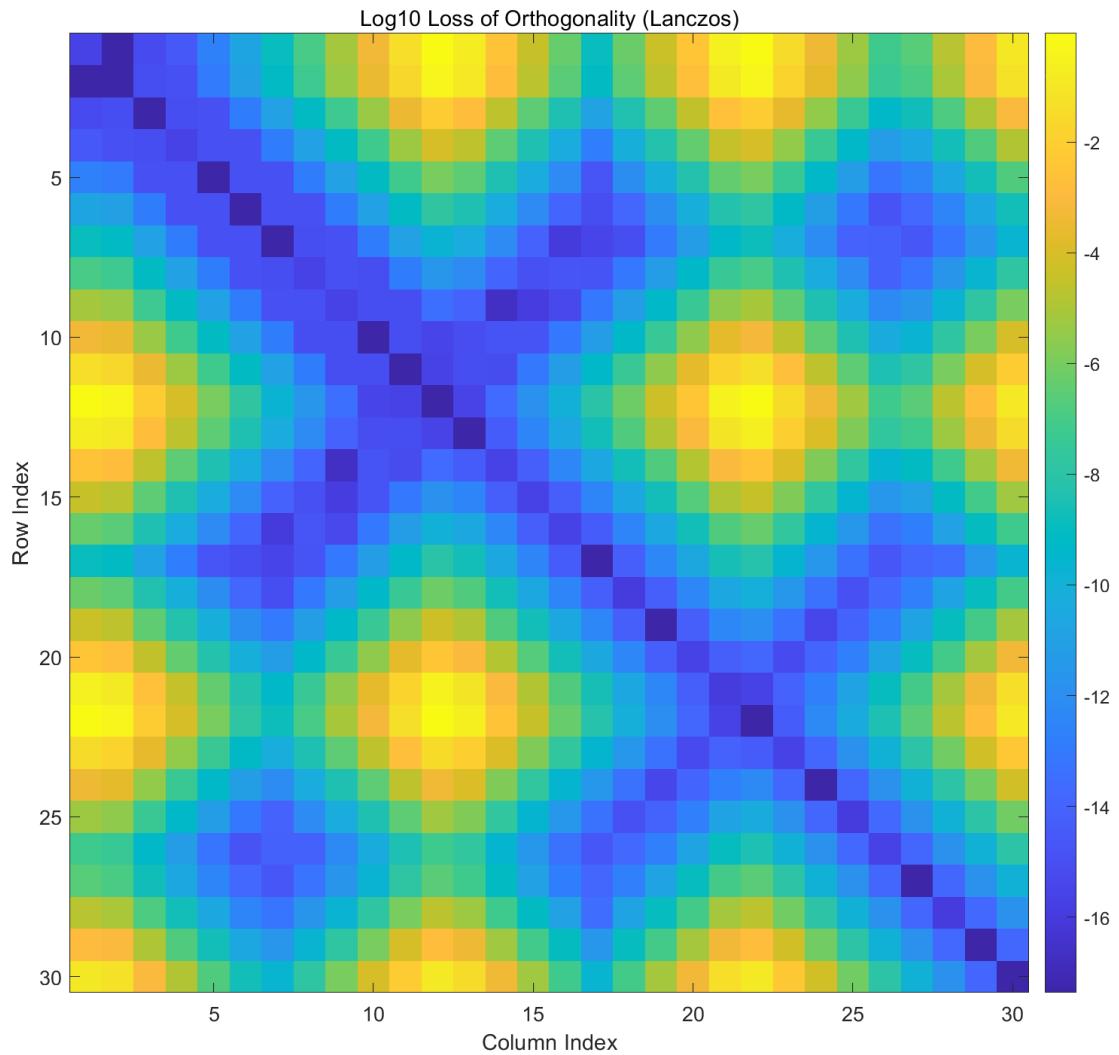
rng(51);
n = 1000;
r = 30;
A = rand(n, n);
A = (A + A') / 2;
b = rand(n, 1);

% Perform the Lanczos algorithm on matrix A with vector b
% r specifies the number of iterations, and 1e-10 is the tolerance for convergence
[Q, T] = Lanczos(A, b, r, 1e-10);

% visualize the loss of orthogonality of the vectors in Q
visualize_orthogonality_loss(Q, 'Log10 Loss of Orthogonality (Lanczos)');

```

可视化正交性损失:



Problem 07 (optional)

Implement the GTH algorithm and verify its componentwise accuracy by some examples.

Solution:

给定实方阵 $A \in \mathbb{R}^{n \times n}$, 其列和均为 1, 即 $\mathbf{1}_n^T A = \mathbf{1}_n^T$

我们的目标是对 A 进行 LU 分解.

首先对 $A^T - I_n$ 进行 Gauss 消元, 得到 $A^T - I_n = LU$

直观起见, 考虑 $n = 4$ 的例子:

(注意到 $A^T - I_n$ 的行和均为 0, 非对角元均为正值, 对角元均为负值)

$$\begin{aligned}
A^T - I_n &= \begin{bmatrix} - & + & + & + \\ + & - & + & + \\ + & + & - & + \\ + & + & + & - \end{bmatrix} \\
&= \begin{bmatrix} 1 & & & \\ -1 & 1 & & \\ - & 1 & 1 & \\ - & & 1 & \end{bmatrix} \begin{bmatrix} - & + & + & + \\ - & + & + & + \\ + & - & + & + \\ + & + & + & + \end{bmatrix} \\
&= \begin{bmatrix} 1 & & & \\ -1 & 1 & & \\ - & 1 & 1 & \\ - & & 1 & \end{bmatrix} \begin{bmatrix} 1 & & & \\ & 1 & & \\ & -1 & 1 & \\ & - & 1 & \end{bmatrix} \begin{bmatrix} - & + & + & + \\ - & + & + & + \\ - & + & - & + \\ + & - & - & + \end{bmatrix} \\
&= \begin{bmatrix} 1 & & & \\ -1 & 1 & & \\ - & 1 & 1 & \\ - & & 1 & \end{bmatrix} \begin{bmatrix} 1 & & & \\ & 1 & & \\ & -1 & 1 & \\ & - & 1 & \end{bmatrix} \begin{bmatrix} 1 & & & \\ & 1 & & \\ & & 1 & \\ & & & 1 \end{bmatrix} \begin{bmatrix} - & + & + & + \\ - & + & + & + \\ - & + & - & + \\ + & - & - & + \end{bmatrix} \\
&= \begin{bmatrix} 1 & & & \\ -1 & 1 & & \\ - & 1 & 1 & \\ - & & 1 & \end{bmatrix} \begin{bmatrix} 1 & & & \\ & 1 & & \\ & -1 & 1 & \\ & - & 1 & \end{bmatrix} \begin{bmatrix} 1 & & & \\ & 1 & & \\ & & 1 & \\ & & & 1 \end{bmatrix} \begin{bmatrix} - & + & + & + \\ - & + & + & + \\ - & + & - & + \\ 0 & - & - & + \end{bmatrix} \\
&= L_3^{-1} L_2^{-1} L_1^{-1} U \\
&= LU
\end{aligned}$$

(消元结束时, 上三角阵最后一行全为零, 这是显然的, 因为行和为零)

我们发现只有对角元位置会发生相消, 而且这个相消是可以避免的.

注意到 $A^T - I_n$ 的行和均为 0, 且列 Gauss 消元会保留 "行和均为 0" 的性质.

因此我们每步 Gauss 消元结束后都要修正对角元,

即将同一行的正元素加起来再取负, 便得到对角元的更精确的计算结果.

上述算法的 Matlab 代码为:

```

function [L, U] = GTH(A)
    % GTH performs Gaussian elimination to decompose a matrix A into lower
    % triangular matrix L and upper triangular matrix U.
    %
    % Input:
    %   A - A square matrix to be decomposed.
    %
    % Output:
    %   L - Lower triangular matrix with ones on the diagonal.
    %   U - Upper triangular matrix.

    % Get the size of matrix A
    [n, ~] = size(A);

    % Perform Gaussian elimination
    for i = 1:n-1

        % Normalize the column below the pivot element
        A(i+1:n, i) = A(i+1:n, i) / A(i, i);

        % Update the submatrix to eliminate elements below the pivot
        A(i+1:n, i+1:n) = A(i+1:n, i+1:n) - A(i+1:n, i) * A(i, i+1:n);

        % Computation for the diagonal elements
        for k = i+1:n
            % Update the diagonal element in the k-th row
            A(k, k) = 0;
            A(k, k) = - sum(A(k, i+1:n));
        end
    end

```

```

    % Output every 100 iterations
    if mod(i, 100) == 0
        fprintf('Iteration %d completed\n', i);
    end

end

% Extract L and U
L = tril(A, -1) + eye(n);
U = triu(A);

end

```

得到 $A^T - I_n = LU$ 后, 我们有:

(注意到 L 是单位下三角阵, 因而是满秩的)

$$\begin{aligned}
 A^T - I_n &= LU \\
 &\Leftrightarrow \\
 A - I_n &= U^T L^T \\
 &\Leftrightarrow \\
 A = (U^T + L^{-T})L^T &= \tilde{L}\tilde{U} \text{ where } \begin{cases} \tilde{L} := U^T + L^{-T} \\ \tilde{U} := L^T \end{cases}
 \end{aligned}$$

这样我们就得到了 A 的 LU 分解 $A = \tilde{L}\tilde{U}$, 可以验证其逐元素绝对误差接近机器精度.

函数调用:

```

rng(51);
n = 2000;
A = rand(n, n);

% Normalize each row of the matrix A using L1 norm
for j = 1:n
    A(:, j) = A(:, j) / norm(A(:, j), 1);
end

% Perform Gaussian elimination using the GTH function
[U_GTH, L_GTH] = GTH(A' - eye(n));
U_GTH = U_GTH';
L_GTH = L_GTH' + U_GTH \ eye(n);

% Calculate the difference between the reconstructed matrix and A
difference_GTH = L_GTH * U_GTH - A;

% Display the maximum absolute error from the reconstruction
disp("Maximum pointwise absolute residual:")
disp(max(abs(difference_GTH(:))));

```

运行结果:

```

Maximum pointwise absolute residual:
8.9214e-16

```