

# 数值算法 Homework 08

Due: Nov. 12, 2024

姓名: 雍崔扬

学号: 21307140051

## Problem 1

Let  $A \in \mathbb{C}^{n \times n}$ ,  $x \in \mathbb{C}^n$

Suppose that  $X = [x, Ax, \dots, A^{n-1}x] \in \mathbb{C}^{n \times n}$  is nonsingular.

Show that  $X^{-1}AX$  is upper Hessenberg.

**Proof:**

给定  $A \in \mathbb{C}^{n \times n}$  和  $x \in \mathbb{C}^n$ , 我们便可以计算序列:

$$\begin{aligned} x_1 &= x \\ x_2 &= Ax_1 = Ax \\ x_3 &= Ax_2 = A^2x \\ &\vdots \\ x_n &= Ax_{n-1} = A^{n-1}x \end{aligned}$$

我们只需计算至  $x_n = A^{n-1}x$  即可

这是因为 Cayley-Hamilton 定理保证了  $A^k$  ( $k \geq n$ ) 可以表示为  $I_n, A, \dots, A^{n-1}$  的线性组合.

(Cayley-Hamilton 定理, Matrix Analysis 定理 2.4.3.2)

设  $p_A(t) := \det(tI_n - A)$  是  $A \in \mathbb{C}^{n \times n}$  的特征多项式, 则我们有  $p_A(A) = 0_{n \times n}$  成立.

换言之, 任意复方阵都满足其特征方程.

我们记:

$$\begin{aligned} X &:= [x_1, x_2, \dots, x_{n-1}, x_n] = [x, Ax, \dots, A^{n-1}x] \\ c &:= -X^{-1}A^n x_1 = [c_1, \dots, c_n]^T \\ H &:= [e_2, e_3, \dots, e_n, -c] = \begin{bmatrix} 0 & & -c_1 \\ 1 & 0 & -c_2 \\ & \ddots & \vdots \\ 1 & & -c_{n-1} \\ \ddots & 0 & -c_n \\ & 1 & -c_n \end{bmatrix} \end{aligned}$$

(其中我们假设  $X$  非奇异)

则我们有:

$$\begin{aligned} AX &= [Ax_1, Ax_2, \dots, Ax_{n-1}, Ax_n] \\ &= [x_2, x_3, \dots, x_n, A^n x_1] \\ &= [Xe_2, Xe_3, \dots, Xe_n, X(-X^{-1}A^n x_1)] \quad (\text{note that } c := X^{-1}A^n x_1) \\ &= X[e_2, e_3, \dots, e_n, -c] \\ &= XH \end{aligned}$$

于是我们有  $X^{-1}AX = H$ , 而  $H$  是一个上 Hessenberg 矩阵(具体来说是一个 Frobenius 酷型)

## Problem 2

Let  $A \in \mathbb{C}^{n \times n}$  be an unreduced upper Hessenberg matrix.

(By "unreduced", we mean  $A_{i+1,i} \neq 0$  ( $\forall 1 \leq i \leq n-1$ ))

Suppose that  $A$  is singular.

Show that the zero eigenvalue appears at the bottom right corner after one QR sweep.

What happens if  $A$  is a singular upper Hessenberg matrix with some  $A_{i+1,i} = 0$ ?

**Solution:**

首先我们注意到奇异的不可约上 Hessenberg 矩阵  $A \in \mathbb{C}^{n \times n}$  的秩  $\text{rank}(A)$  一定为  $n-1$

因此  $A$  的 QR 分解形如:

$$A = QR$$
$$\begin{cases} Q = G_{n-1,n} \cdots G_{1,2} \in \mathbb{C}^{n \times n} \text{ where } G_{1,2}, \dots, G_{n-1,n} \text{ are Givens matrices} \\ R = \begin{bmatrix} R_1 & r \\ 0 & 0 \end{bmatrix} \in \mathbb{C}^{n \times n} \text{ where } R_1 \in \mathbb{C}^{(n-1) \times (n-1)} \text{ is a non-singular upper triangle matrix} \end{cases}$$

于是我们在计算  $\tilde{A} = RQ = RG_{1,2}^H \cdots G_{n-1,n}^H$  时, 其第  $n$  行的所有元素必然全为零.

这是因为  $R$  的第  $n$  行的所有元素全为零, 所以无论怎么做列线性组合, 第  $n$  行的所有元素必然全为零.

因此一次 QR 迭代得到的  $\tilde{A} = RQ = Q^H A Q$  的零特征值将位于  $(n, n)$  位置.

以  $n = 4$  的情况为例,  $A$  的 Givens QR 分解过程如下:

$$A = \begin{bmatrix} * & * & * & * \\ * & * & * & * \\ * & * & * & * \\ * & * & * & * \end{bmatrix}$$

$$G_{1,2}A = \begin{bmatrix} * & * & * & * \\ 0 & * & * & * \\ * & * & * & * \\ * & * & * & * \end{bmatrix}$$

$$G_{2,3}(G_{1,2}A) = \begin{bmatrix} * & * & * & * \\ 0 & * & * & * \\ 0 & * & * & * \\ * & * & * & * \end{bmatrix}$$

$$G_{3,4}(G_{2,3}G_{1,2}A) = \begin{bmatrix} * & * & * & * \\ 0 & * & * & * \\ 0 & * & * & * \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

下面我们计算  $\tilde{A} = RQ$ :

$$R = \begin{bmatrix} * & * & * & * \\ 0 & * & * & * \\ 0 & * & * & * \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$RG_{1,2}^H = \begin{bmatrix} * & * & * & * \\ * & * & * & * \\ 0 & * & * & * \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$(RG_{1,2}^H)G_{2,3}^H = \begin{bmatrix} * & * & * & * \\ * & * & * & * \\ * & * & * & * \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$(RG_{1,2}^H G_{2,3}^H)G_{3,4}^H = \begin{bmatrix} * & * & * & * \\ * & * & * & * \\ * & * & * & * \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

What happens if  $A$  is a singular upper Hessenberg matrix with some  $A_{i+1,i} = 0$ ?

那么一步 QR 后, 零特征值会落到第一个不可约 Hessenberg 矩阵的底部.

(如果上三角阵  $R$  可以是阶梯形式的话, 那么也可以把零特征值落到整个 Hessenberg 矩阵的底部)

### Problem 3

Let:

$$A = \begin{bmatrix} 0 & & & 1 \\ 1 & 0 & & \\ \ddots & \ddots & \ddots & \\ & 1 & 0 & \\ & & 1 & 0 \end{bmatrix}$$

What can you say about the convergence of:

- ① the naive QR algorithm
- ② Francis' double-shift QR algorithm

### (0) Helper Functions

计算 Givens 变换的算法为:

```

function [c, s] = Givens(a, b)
    if b == 0
        c = 1; s = 0
    else
        if |b| > |a|
            t =  $\frac{a}{b}$ ; s =  $\frac{1}{\sqrt{1+t^2}}$ ; c = st
        else
            t =  $\frac{b}{a}$ ; c =  $\frac{1}{\sqrt{1+t^2}}$ ; s = ct
        end
    end
end

```

其 Matlab 代码如下:

```

function [c, s] = Givens(a, b)
    % Givens 旋转, 计算 cos 和 sin
    if b == 0
        c = 1;
        s = 0;
    else
        if abs(b) > abs(a)
            t = a / b;
            s = 1 / sqrt(1 + t^2);
            c = s * t;
        else
            t = b / a;
            c = 1 / sqrt(1 + t^2);
            s = c * t;
        end
    end
end

```

复数域上的 Householder 变换已于 Homework 02 Problem 02 给出:

```

function [v, beta] = Complex_Householder(x)
    % This function computes the Householder vector 'v' and scalar 'beta' for
    % a given complex vector 'x'. This transformation is used to create zeros
    % below the first element of 'x' by reflecting 'x' along a specific direction.

    n = length(x);
    x = x / norm(x, inf); % Normalize x by its infinity norm to avoid numerical issues

    % Copy all elements of 'x' except the first into 'v'
    v = zeros(n, 1);
    v(2:n) = x(2:n);

    % Compute sigma as the squared 2-norm of the elements of x starting from the second element
    sigma = norm(x(2:n), 2)^2;

    % Check if sigma is near zero, which would mean 'x' is already close to a scalar multiple of e_1
    if sigma < 1e-10
        beta = 0; % If sigma is close to zero, set beta to zero (no transformation needed)
    else
        % Determine gamma to account for the argument of complex number x(1)
        if abs(x(1)) < 1e-10
            gamma = 1; % If x(1) is close to zero, set gamma to 1
        else
            gamma = x(1) / abs(x(1)); % Otherwise, set gamma to x(1) divided by its magnitude
        end

        % Compute alpha as the Euclidean norm of x, including x(1) and sigma
        alpha = sqrt(abs(x(1))^2 + sigma);

        % Compute the first element of 'v' to avoid numerical cancellation
        v(1) = -gamma * sigma / (abs(x(1)) + alpha);

        % Calculate 'beta', the scaling factor of the Householder transformation
        beta = 2 * abs(v(1))^2 / (abs(v(1))^2 + sigma);

        % Normalize the vector 'v' by v(1) to ensure that the first element is 1,
        % allowing for simplified storage and computation of the transformation
        v = v / v(1);
    end
end

```

## (1) Algorithm 1

上 Hessenberg 矩阵的显式 QR 迭代:

---

Given Hessenberg matrix  $H \in \mathbb{R}^{n \times n}$

---


$$Q = I_n$$

$$G = \text{Zeros}(n - 1, 2)$$

$$\text{for } k = 1 : n - 1$$

$$[c, s] = \text{Givens}(H(k, k), H(k + 1, k))$$

$$G(k, 1 : 2) = [c, s]$$

$$H(k : k + 1, k : n) = \begin{bmatrix} c & s \\ -s & c \end{bmatrix} H(k : k + 1, k : n)$$

$$Q(1 : n, k : k + 1) = Q(1 : n, k : k + 1) \begin{bmatrix} c & s \\ -s & c \end{bmatrix}^T$$

$$\text{end}$$

$$R = H$$


---


$$\text{for } k = 1 : n - 1$$

$$[c, s] = G(k, 1 : 2)$$

$$H(1 : k + 1, k : k + 1) = H(1 : k + 1, k : k + 1) \begin{bmatrix} c & s \\ -s & c \end{bmatrix}^T$$

$$\text{end}$$

$$\tilde{H} = H$$

Matlab 代码为:

```

function [Q, R, H_tilde] = Hessenberg_Givens_Reduction(H)
    % Given Hessenberg matrix H, apply Givens rotations to reduce it.
    % Outputs:
    %   Q       : Orthogonal matrix obtained from applying Givens rotations.
    %   R       : Resulting Hessenberg matrix after reduction.
    %   H_tilde : Modified Hessenberg matrix after applying additional Givens rotations.

    % Initialize the size of the matrix
    n = size(H, 1);

    % Step 1: Initialize Q as the identity matrix, and G to store Givens rotations
    Q = eye(n);
    G = zeros(n-1, 2); % Each row will store [c, s] values for each rotation

    % Step 2: Apply Givens rotations to H
    for k = 1:n-1
        % Compute Givens rotation parameters [c, s] for H(k, k) and H(k, k+1)
        [c, s] = Givens(H(k, k), H(k+1, k));

        % Store the Givens rotation parameters in G
        G(k, :) = [c, s];

        % Apply Givens rotation to H in rows k and k+1 from column k to end
        H(k:k+1, k:n) = [c, s; -s, c] * H(k:k+1, k:n);

        % Apply the transpose of the Givens rotation to Q in columns k and k+1
        Q(:, k:k+1) = Q(:, k:k+1) * [c, s; -s, c]';

    end

    % Store the result of H as R after the first Givens reduction phase
    R = H;

    % Step 3: Apply stored Givens rotations to update H
    for k = 1:n-1
        % Retrieve the Givens rotation parameters [c, s] from G
        c = G(k, 1);
        s = G(k, 2);

        % Apply the transpose of the Givens rotation to H in rows 1:k+1 and columns k:k+1
        H(1:k+1, k:k+1) = H(1:k+1, k:k+1) * [c, s; -s, c]';

    end

    % Store the modified Hessenberg matrix after the second phase as H_tilde
    H_tilde = H;
end

```

应用上述算法我们发现每步得到的正交矩阵  $Q_k$  均为:

$$Q_k \equiv Q := \begin{bmatrix} 0 & & & -1 \\ 1 & 0 & & \\ & \ddots & \ddots & \\ & & 1 & 0 \\ & & & 1 & 0 \end{bmatrix} (\forall k \in \mathbb{Z}_+)$$

可以验证  $A_n = (Q^n)^T A (Q^n) = A$ , 因此对上 Hessenberg 矩阵  $A$  的显式 QR 迭代不收敛.  
以  $n = 4$  为例:

```

n = 4;
% Set the subdiagonal to 1
A = diag(ones(n-1, 1), -1);

% Set the top right element to 1
A(1, n) = 1;

[Q1, R1, A1] = Hessenberg_Givens_Reduction(A);
[Q2, R2, A2] = Hessenberg_Givens_Reduction(A1);
[Q3, R3, A3] = Hessenberg_Givens_Reduction(A2);
[Q4, R4, A4] = Hessenberg_Givens_Reduction(A3);

% Display results in a formatted way
disp("Results of Hessenberg_Givens_Reduction:")

% Display Q1, R1, A1 side by side
disp("Round 1: [Q1, R1, A1]")
disp([Q1, R1, A1])

% Display Q2, R2, A2 side by side
disp("Round 2: [Q2, R2, A2]")
disp([Q2, R2, A2])

% Display Q3, R3, A3 side by side
disp("Round 3: [Q3, R3, A3]")
disp([Q3, R3, A3])

% Display Q4, R4, A4 side by side
disp("Round 4: [Q4, R4, A4]")
disp([Q4, R4, A4])

```

运行结果:

```

Results of Hessenberg_Givens_Reduction:
Round 1: [Q1, R1, A1]
0 0 0 -1 1 0 0 0 0 0 0 -1
1 0 0 0 0 1 0 0 1 0 0 0
0 1 0 0 0 0 1 0 0 1 0 0
0 0 1 0 0 0 0 0 -1 0 0 -1 0

Round 2: [Q2, R2, A2]
0 0 0 -1 1 0 0 0 0 0 0 -1
1 0 0 0 0 1 0 0 1 0 0 0
0 1 0 0 0 0 -1 0 0 -1 0 0
0 0 1 0 0 0 0 1 0 0 1 0

Round 3: [Q3, R3, A3]
0 0 0 -1 1 0 0 0 0 0 0 -1
1 0 0 0 0 -1 0 0 -1 0 0 0
0 1 0 0 0 0 1 0 0 1 0 0
0 0 1 0 0 0 0 1 0 0 1 0

Round 4: [Q4, R4, A4]
0 0 0 -1 -1 0 0 0 0 0 0 1
1 0 0 0 0 1 0 0 1 0 0 0
0 1 0 0 0 0 1 0 0 1 0 0
0 0 1 0 0 0 0 1 0 0 1 0

```

## (2) Algorithm 2

(Francis 双位移的 QR 迭代算法, 数值线性代数, 算法 6.4.2)

Given Hessenberg matrix  $H \in \mathbb{R}^{n \times n}$

$$t = \text{tr} \begin{pmatrix} h_{n-1,n-1} & h_{n-1,n} \\ h_{n,n-1} & h_{n,n} \end{pmatrix} = h_{n-1,n-1} + h_{n,n}$$

$$s = \det \begin{pmatrix} h_{n-1,n-1} & h_{n-1,n} \\ h_{n,n-1} & h_{n,n} \end{pmatrix} = h_{n-1,n-1}h_{n,n} - h_{n-1,n}h_{n,n-1}$$

$$\begin{cases} m_{11} = h_{11}^2 + h_{12}h_{21} - th_{11} + s \\ m_{21} = h_{21}(h_{11} + h_{22} - t) \\ m_{31} = h_{21}h_{32} \end{cases} \quad (\text{note that } Me_1 = \begin{bmatrix} m_{11} \\ m_{21} \\ m_{31} \\ 0 \\ \vdots \\ 0 \end{bmatrix} = \begin{bmatrix} h_{11}^2 + h_{12}h_{21} - th_{11} + s \\ h_{21}(h_{11} + h_{22} - t) \\ h_{21}h_{32} \\ 0 \\ \vdots \\ 0 \end{bmatrix} \text{ where } M = H^2 - tH + sI)$$

$$[v, \beta] = \text{Householder} \begin{pmatrix} m_{11} \\ m_{21} \\ m_{31} \end{pmatrix} \quad (\text{case of } k=0)$$

$$H(1:3, 1:n) = (I_3 - \beta vv^T)H(1:3, 1:n) = H(1:3, 1:n) - (\beta v)(v^T H(1:3, 1:n))$$

$$H(1:4, 1:3) = H(1:4, 1:3)(I_3 - \beta vv^T) = H(1:4, 1:3) - (H(1:4, 1:3)v)(\beta v)^T$$

for  $k = 1:n-4$

$$[v, \beta] = \text{Householder}(H(k+1:k+3, k))$$

$$H(k+1:k+3, k:n) = (I_3 - \beta vv^T)H(k+1:k+3, k:n) = H(k+1:k+3, k:n) - (\beta v)(v^T H(k+1:k+3, k:n))$$

$$H(1:k+4, k+1:k+3) = H(1:k+4, k+1:k+3)(I_3 - \beta vv^T) = H(1:k+4, k+1:k+3) - (H(1:k+4, k+1:k+3)v)(\beta v)^T$$

end

$$[v, \beta] = \text{Householder}(H(n-2:n, n-3)) \quad (\text{case of } k=n-3)$$

$$H(n-2:n, n-3:n) = (I_3 - \beta vv^T)H(n-2:n, n-3:n) = H(n-2:n, n-3:n) - (\beta v)(v^T H(n-2:n, n-3:n))$$

$$H(1:n, n-2:n) = H(1:n, n-2:n)(I_3 - \beta vv^T) = H(1:n, n-2:n) - (H(1:n, n-2:n)v)(\beta v)^T$$

$$[v, \beta] = \text{Householder}(H(n-1:n, n-2)) \quad (\text{case of } k=n-2)$$

$$H(n-1:n, n-2:n) = (I_2 - \beta vv^T)H(n-1:n, n-2:n) = H(n-1:n, n-2:n) - (\beta v)(v^T H(n-1:n, n-2:n))$$

$$H(1:n, n-1:n) = H(1:n, n-1:n)(I_2 - \beta vv^T) = H(1:n, n-1:n) - (H(1:n, n-1:n)v)(\beta v)^T$$

其 Matlab 代码为: (加入了调试信息)

```

function H_tilde = Francis_Double_Shift_QR_Iteration(H)
    % Given a Hessenberg matrix H, apply Francis Double Shift QR Iteration
    % and return the modified Hessenberg matrix H_tilde.

    n = size(H, 1); % Dimension of H
    disp("Original H:")
    disp(H);

    % Step 1: Compute the trace and determinant for the 2x2 bottom right submatrix
    t = H(n-1, n-1) + H(n, n); % Trace of bottom-right 2x2 block
    s = H(n-1, n-1) * H(n, n) - H(n-1, n) * H(n, n-1); % Determinant

    % Step 2: Define the components of Me_1 vector
    m11 = H(1, 1)^2 + H(1, 2) * H(2, 1) - t * H(1, 1) + s;
    m21 = H(2, 1) * (H(1, 1) + H(2, 2) - t);
    m31 = H(2, 1) * H(3, 2);
    Me1 = [m11; m21; m31];

    % Step 3: Apply Householder transformation to Me1 (initial case k=0)
    [v, beta] = Complex_Householder(Me1); % Compute Householder vector and scalar
    % Apply transformation to H
    H(1:3, 1:n) = H(1:3, 1:n) - beta * (v * (v' * H(1:3, 1:n)));
    H(1:4, 1:3) = H(1:4, 1:3) - (H(1:4, 1:3) * v) * (beta * v)';
    disp("After applying Householder of Me1:")
    disp(H);

    % Step 4: Iterate for k = 1 to n-4
    for k = 1:n-4
        [v, beta] = Complex_Householder(H(k+1:k+3, k)); % Compute Householder for each block
        H(k+1:k+3, k:n) = H(k+1:k+3, k:n) - beta * (v * (v' * H(k+1:k+3, k:n)));
        H(1:k+4, k+1:k+3) = H(1:k+4, k+1:k+3) - (H(1:k+4, k+1:k+3) * v) * (beta * v)';
        fprintf("After applying the %d-th Householder:\n", k);
        disp(H);
    end

    % Step 5: Handle case for k = n-3
    [v, beta] = Complex_Householder(H(n-2:n, n-3));
    H(n-2:n, n-3:n) = H(n-2:n, n-3:n) - beta * (v * (v' * H(n-2:n, n-3:n)));
    H(1:n, n-2:n) = H(1:n, n-2:n) - (H(1:n, n-2:n) * v) * (beta * v)';
    fprintf("After applying the %d-th Householder:\n", n-3);
    disp(H);

    % Step 6: Handle case for k = n-2

```

```

[v, beta] = Complex_Householder(H(n-1:n, n-2));
H(n-1:n, n-2:n) = H(n-1:n, n-2:n) - beta * (v * (v' * H(n-1:n, n-2:n)));
H(1:n, n-1:n) = H(1:n, n-1:n) - (H(1:n, n-1:n) * v) * (beta * v)';
fprintf("After applying the %d-th Householder: (Final matrix)\n", n-2);
disp(H);

% Output the modified Hessenberg matrix
H_tilde = H;
end

```

函数调用:

```

n = 5;
% Set the subdiagonal to 1
A = diag(ones(n-1, 1), -1);

% Set the top right element to 1
A(1, n) = 1;

A1 = Francis_Double_Shift_QR_Iteration(A);

```

运行结果:

```

Original H:
 0   0   0   0   1
 1   0   0   0   0
 0   1   0   0   0
 0   0   1   0   0
 0   0   0   1   0

After applying Householder of Mel:
 0   1   0   0   0
 0   0   1   0   0
 0   0   0   0   1
 1   0   0   0   0
 0   0   0   1   0

After applying the 1-th Householder:
 0   0   0   1   0
 1   0   0   0   0
 0   0   0   0   1
 0   0   1   0   0
 0   1   0   0   0

After applying the 2-th Householder:
 0   0   0   1   0
 1   0   0   0   0
 0   1   0   0   0
 0   0   0   0   1
 0   0   1   0   0

After applying the 3-th Householder: (Final matrix)
 0   0   0   0   1
 1   0   0   0   0
 0   1   0   0   0
 0   0   1   0   0
 0   0   0   1   0

```

我们发现经过一轮 Francis 双步位移隐式 QR 迭代，矩阵  $A$  并无变化，因此算法不收敛。

## Problem 4

$$\text{Let } A = \begin{bmatrix} a & d \\ b & c \end{bmatrix} \in \mathbb{R}^{2 \times 2}$$

Design an algorithm to compute an orthogonal matrix  $Q \in \mathbb{R}^{2 \times 2}$  such that  $Q^T A Q = \begin{bmatrix} b & d \\ 0 & a \end{bmatrix}$

(optional) What happens if the matrix  $A$  is complex?

(optional) Write a subprogram to perform diagonal swapping of the real Schur form:

$$A = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}$$

where  $A_{11}$  and  $A_{22}$  are either  $1 \times 1$  or  $2 \times 2$

## Part (1)

Let  $A = \begin{bmatrix} a & d \\ b & c \end{bmatrix} \in \mathbb{R}^{2 \times 2}$

Design an algorithm to compute an orthogonal matrix  $Q \in \mathbb{R}^{2 \times 2}$  such that  $Q^T A Q = \begin{bmatrix} b & d \\ a & c \end{bmatrix}$

**Solution:**

记  $Q = \begin{bmatrix} c & s \\ -s & c \end{bmatrix} \in \mathbb{R}^{2 \times 2}$  (满足  $c^2 + s^2 = 1$ )

于是我们有:

$$\begin{aligned} Q^T A Q &= \begin{bmatrix} c & s \\ -s & c \end{bmatrix}^T \begin{bmatrix} a & d \\ b & c \end{bmatrix} \begin{bmatrix} c & s \\ -s & c \end{bmatrix} \\ &= \begin{bmatrix} c & -s \\ s & c \end{bmatrix} \begin{bmatrix} ac - ds & as + dc \\ -bs & bc \end{bmatrix} \\ &= \begin{bmatrix} ac^2 - dsc + bs^2 & asc + dc^2 - bsc \\ asc - ds^2 - bsc & as^2 + dsc + bc^2 \end{bmatrix} \\ &= \begin{bmatrix} b & d \\ 0 & a \end{bmatrix} \end{aligned}$$

问题归结为求解方程组:

$$\begin{cases} ac^2 - dsc + bs^2 = b \\ asc + dc^2 - bsc = d \\ asc - ds^2 - bsc = 0 \\ as^2 + dsc + bc^2 = a \\ c^2 + s^2 = 1 \end{cases}$$

根据第三个等式我们有  $s[(a - b)c - ds] = 0$

- 若  $s = 0$ , 则  $c = \pm 1$ , 根据第四个等式我们有  $a = b$  (因此这种取法可以应对  $a = b$  的情况)
- 若  $s \neq 0$  (此情况默认  $a \neq b$ ), 则我们有  $(a - b)c - ds = 0$ 
  - 若  $d = 0$ , 则我们可取  $\begin{cases} c = 0 \\ s = 1 \end{cases}$  (这个情况可以合并入下面的情况)
  - 若  $d \neq 0$ , 则我们有  $t = \frac{s}{c} = \frac{a-b}{d}$ , 可取  $\begin{cases} c = \frac{d}{\sqrt{(a-b)^2+d^2}} \\ s = \frac{a-b}{\sqrt{(a-b)^2+d^2}} \end{cases}$

综上所述, 我们得到如下算法:

```
function [c, s] = swap_diagonal(A)
    % Extract elements from matrix A
    a = A(1,1);
    d = A(1,2);
    b = A(2,2);

    % Initialize values of c and s based on conditions
    if abs(a-b) / (abs(a) + abs(b)) < 1e-10
        % Case of a = b
        c = 1;
        s = 0;
    else
        % General case: a ≠ b
        % Calculate c and s based on derived t value
        c = d / sqrt((a - b)^2 + d^2);
        s = (a - b) / sqrt((a - b)^2 + d^2);
    end
end
```

函数调用:

```
rng(51);
A = triu(rand(2, 2));
disp("Original A:");
disp(A);

[c, s] = swap_diagonal(A);
Q = [c, s; -s, c];
A_tilde = Q' * A * Q;
disp("Swapped A:");
disp(A_tilde);
```

运行结果:

```
Original A:
0.6757  0.3433
0       0.6440
```

```
Swapped A:
0.6440  0.3433
0       0.6757
```

## Part (2)

What happens if the matrix  $A$  is complex?

Let  $A = \begin{bmatrix} a & d \\ b & c \end{bmatrix} \in \mathbb{C}^{2 \times 2}$

Design an algorithm to compute an unitary matrix  $Q \in \mathbb{C}^{2 \times 2}$  such that  $Q^H A Q = \begin{bmatrix} b & d \\ a & c \end{bmatrix}$

**Solution:**

记  $Q = \begin{bmatrix} c & s \\ -\bar{s} & c \end{bmatrix} \in \mathbb{C}^{2 \times 2}$  (满足  $\begin{cases} c^2 + |s|^2 = 1 \\ c \in \mathbb{R} \\ s \in \mathbb{C} \end{cases}$ )

于是我们有:

$$\begin{aligned} Q^H A Q &= \begin{bmatrix} c & s \\ -\bar{s} & c \end{bmatrix}^H \begin{bmatrix} a & d \\ b & c \end{bmatrix} \begin{bmatrix} c & s \\ -\bar{s} & c \end{bmatrix} \\ &= \begin{bmatrix} c & -\bar{s} \\ \bar{s} & c \end{bmatrix} \begin{bmatrix} ac - d\bar{s} & as + dc \\ -b\bar{s} & bc \end{bmatrix} \\ &= \begin{bmatrix} ac^2 - d\bar{s}c + b|s|^2 & asc + dc^2 - bsc \\ a\bar{s}c - d\bar{s}^2 - b\bar{s}c & a|s|^2 + d\bar{s}c + bc^2 \end{bmatrix} \\ &= \begin{bmatrix} b & d \\ 0 & a \end{bmatrix} \end{aligned}$$

问题归结为求解方程组:

$$\begin{cases} ac^2 - d\bar{s}c + b|s|^2 = b \\ asc + dc^2 - bsc = d \\ a\bar{s}c - d\bar{s}^2 - b\bar{s}c = 0 \\ a|s|^2 + d\bar{s}c + bc^2 = a \\ c^2 + s^2 = 1 \end{cases}$$

根据第三个等式我们有  $\bar{s}[(a - b)c - d\bar{s}] = 0$

- 若  $\bar{s} = 0$ , 则  $c = \pm 1$ , 根据第四个等式我们有  $a = b$  (因此这种取法可以应对  $a = b$  的情况)
- 若  $\bar{s} \neq 0$  (此情况默认  $a \neq b$ ), 则我们有  $(a - b)c - d\bar{s} = 0$

◦ 若  $d = 0$ , 则我们可取  $\begin{cases} c = 0 \\ s = 1 \end{cases}$  (这个情况可以合并入下面的情况)

◦ 若  $d \neq 0$ , 则我们有  $t = \frac{\bar{s}}{c} = \frac{a-b}{d} = \frac{(a-b)\bar{d}}{|d|^2}$ , 可取  $\begin{cases} c = \frac{|d|^2}{\sqrt{|a-b|^2|d|^2+|d|^4}} = \frac{|d|}{\sqrt{|a-b|+|d|^2}} \\ \bar{s} = \frac{(a-b)\bar{d}}{|d|\sqrt{(a-b)^2+d^2}} \\ s = \frac{(\bar{a}-\bar{b})d}{|d|\sqrt{(a-b)^2+d^2}} \end{cases}$

(可以放心, Wilkinson 模型保证两个浮点数  $a - b$  的运算不会相消)

综上所述, 我们得到如下算法:

```
function [c, s] = swap_complex(A)
    % Extract elements from the Hermitian matrix A
    a = A(1,1);
    d = A(1,2);
    b = A(2,2);

    % Initialize values of c and s based on conditions
    if abs(a - b) < 1e-10 * (abs(a) + abs(b))
        % Case of a = b (small difference, considering floating point precision)
        c = 1;      % c = 1 for identity-like rotation
        s = 0;      % s = 0 implies no rotation needed
    else
        % General case: a != b
        % Calculate the magnitude of d and the rotation terms
        norm_factor = sqrt(abs(a - b)^2 + abs(d)^2);

        % Compute c and s
        c = abs(d) / norm_factor;
        s = (conj(a) - conj(b)) * d / (abs(d) * norm_factor); % s is a complex number
    end
```

```
end
```

函数调用:

```

rng(51);
A = triu(rand(2, 2) + 1i * rand(2, 2));
disp("Original A:");
disp(A);

[c, s] = swap_compTex(A);
Q = [c, s; -conj(s), c];
A_tilde = Q' * A * Q;
disp("Swapped A:");
disp(A_tilde);

```

运行结果:

```

Original A:
0.6757 + 0.2842i  0.3433 + 0.1577i
0.0000 + 0.0000i  0.6440 + 0.3880i

Swapped A:
0.6440 + 0.3880i  0.3433 + 0.1577i
0.0000 + 0.0000i  0.6757 + 0.2842i

```

## Part (3)

Write a subprogram to perform diagonal swapping of the real Schur form:

$$A = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}$$

where  $A_{11}$  and  $A_{22}$  are either  $1 \times 1$  or  $2 \times 2$

(参考论文: On Swapping Diagonal Blocks in Real Schur Form (Z. Bai & J. Demmel))

**Solution:**

设  $A_{11} \in \mathbb{R}^{p \times p}$ ,  $A_{22} \in \mathbb{R}^{q \times q}$

- ① 使用全选主元 Gauss 消去法求解:

$$A_{11}X - XA_{22} = \gamma A_{12}$$

$\Leftrightarrow$

$$(I_q \otimes A_{11} - A_{22}^T \otimes I_p) \text{vec}(X) = \gamma \text{vec}(A_{12})$$

其中我们假设  $A_{11}$  和  $A_{22}$  没有公共特征值 (这样 Sylvester 定理保证了上述方程组有唯一解)  
而  $\gamma \leq 1$  是一个预设的防止上溢的缩放因子.

全选主元 Gauss 消去法过程中, 若有对角元非常小, 则设置其为  $\text{eps} \cdot \|I_q \otimes A_{11} - A_{22}^T \otimes I_p\|_F$

- ② 使用 Householder 变换计算  $\begin{bmatrix} -X \\ \gamma I_q \end{bmatrix}$  的 QR 分解

- ③ 计算  $Q^T A Q$ :

$$Q^T A Q = \begin{bmatrix} \tilde{A}_{22} & \tilde{A}_{12} \\ \tilde{A}_{21} & \tilde{A}_{11} \end{bmatrix}$$

可以证明  $\tilde{A}_{11}$  和  $A_{11}$  具有相同的特征值, 而  $\tilde{A}_{22}$  和  $A_{22}$  具有相同的特征值.

- ④ 若  $\|\tilde{A}_{21}\|_F \leq \text{eps} \cdot \|A\|_F$ , 则我们记录  $\tilde{A} = Q^T A Q$  (并将  $\tilde{A}_{21}$  设为  $0_{p \times q}$ )  
否则报错并退出 (这表明向后不稳定, 结果不被接受)
- ⑤ 标准化  $2 \times 2$  对角块 (如果有的话)

**粗略的 Matlab 实现:**

```

function A_swapped = diagonal_swap(A, gamma)
    % This function performs the diagonal swapping of the real schur form
    % A is the matrix in real Schur form (block diagonal matrix)
    % gamma is a scaling factor to prevent overflow

    % Step 1: Determine p and q based on the size of A
    n = size(A, 1);
    if n == 2
        p = 1;
        q = 1;
    elseif n == 4
        p = 2;
    end

```

```

q = 2;
elseif n == 3
    % If n = 3, check the off-diagonal elements to determine p and q
    if A(2,1) == 0 % If A12 is zero, then p = 1, q = 2
        p = 1;
        q = 2;
    else % Otherwise, p = 2, q = 1
        p = 2;
        q = 1;
    end
else
    error('Matrix size n should be 2, 3, or 4.');
end

% Step 2: Extract the block components A11, A12, and A22
A11 = A(1:p, 1:p); % Top-left block
A12 = A(1:p, p+1:end); % Top-right block
A22 = A(p+1:end, p+1:end); % Bottom-right block

% Step 3: Solve the system of equations: A11*X - X*A22 = gamma*A12
% Construct the system (I_q ⊗ A11 - A22^T ⊗ I_p) * vec(X) = gamma * vec(A12)
X_vec = solve_system(A11, A22, A12, gamma);

% Reshape X_vec back to matrix form
X = reshape(X_vec, [p, q]);

% Step 4: Apply Householder transformation
% Form the vector [-X; gamma*I_q]
[Q, ~] = qr([-X; gamma * eye(q)]);

% Step 5: Compute Q^T * A * Q
A_swapped = Q' * A * Q;

% Extract the (2, 1) blocks from A_swapped
A21_tilde = A_swapped(q+1:end, 1:q);

% Step 6: Check if the off-diagonal block A12_tilde is small enough
if norm(A21_tilde, 'fro') <= eps * norm(A_swapped, 'fro')
    % Set the off-diagonal block to zero
    A_swapped(q+1:end, 1:q) = zeros(p, q);
else
    error('fatal: the (2,1) block of A_swapped is non-zero!')
end
end

function X_vec = solve_system(A11, A22, A12, gamma)
    % Solve the system (I_q ⊗ A11 - A22^T ⊗ I_p) * vec(X) = gamma * vec(A12)
    % Use a direct solution or an iterative solver depending on the structure
    % of the system (simplified here as a dense solver for illustrative purposes)

    [p, q] = size(A12);

    % Construct the Kronecker product matrix (I_q ⊗ A11 - A22^T ⊗ I_p)
    K = kron(eye(q), A11) - kron(A22', eye(p));

    % Vectorize the right-hand side (gamma * A12)
    rhs = gamma * A12(:);

    % Solve the linear system
    X_vec = K \ rhs; % This uses the backslash operator, which is efficient in MATLAB
end

```

函数调用:

```

A = [1, 1, 1;
      0, 2, 3;
      0,-3, 2];
gamma = 1;
A_swapped = diagonal_swap(A, gamma);

disp("Original A:")
disp(A);
disp(eig(A));

disp("A_swapped:");
disp(A_swapped);
disp(eig(A_swapped));

```

运行结果:

```
original A:  
 1 1 1  
 0 2 3  
 0 -3 2  
  
1.0000 + 0.0000i  
2.0000 + 3.0000i  
2.0000 - 3.0000i  
  
A_swapped:  
 2.2069 3.1919 -0.1017  
-2.8330 1.7931 -1.2999  
 0 0 1.0000  
  
2.0000 + 3.0000i  
2.0000 - 3.0000i  
1.0000 + 0.0000i
```

## Problem 5

Implement the following algorithms for Hessenberg reduction:

- (a) using Householder reflections;
- (b) using Arnoldi process based on modified Gram-Schmidt orthogonalization.

Randomly generate a few matrices and compute the corresponding Hessenberg decomposition  $A = QHQ^H$

Check the accuracy in terms of  $\|Q^H A Q - H\|_F$  and  $\|Q^H Q - I\|_F$  for your Hessenberg reduction implementations.

What do you observe?

(optional) Perturb the matrix  $A$  a little bit. How do  $Q$  and  $H$  change accordingly?

### Part (1)

Implement Hessenberg reduction for  $A \in \mathbb{C}^{n \times n}$  by Householder reflections.

Check the accuracy in terms of  $\|Q^H A Q - H\|_F$  and  $\|Q^H Q - I\|_F$  for your Hessenberg reduction implementations.

Perturb the matrix  $A$  a little bit.

How do  $Q$  and  $H$  change accordingly?

#### Solution:

使用 Householder 变换进行上 Hessenberg 化的算法:

```
Given  $A \in \mathbb{C}^{n \times n}$   
_____  
 $W = 0_{n-1, n-2}$   
 $Y = 0_{n-1, n-2}$   
 $b = 0_{n-2}$   
for  $k = 1 : n - 2$   
   $[v, \beta] = \text{Householder}(A(k+1:n, k))$   
   $A(k+1:n, k:n) = (I_{n-k} - \beta vv^H)A(k+1:n, k:n) = A(k+1:n, k:n) - (\beta v)(v^H A(k+1:n, k:n))$   
   $A(1:n, k+1:n) = A(1:n, k+1:n)(I_{n-k} - \beta vv^H) = A(1:n, k+1:n) - (A(1:n, k+1:n)v)(\beta v)^H$   
   $Y(k:n-1, k) = v$   
   $b(k) = \beta$   
end  
for  $k = 1 : n - 2$   
   $v = Y(k:n-1, k)$   
   $\beta = b(k)$   
  if  $k = 1$   
     $W(1:n-1, 1) = -\beta v$   
  else  
     $W(1:n-1, k) = W(1:n-1, 1:k-1)[Y(k:n-1, 1:k-1)^H v]$   
     $W(k:n-1, k) = W(k:n-1, k) + v$   
     $W(1:n-1, k) = -\beta W(1:n-1, k)$   
  end  
end  
 $H = A$   
 $Q = \begin{bmatrix} 1 & \\ & I_{n-1} + WY^H \end{bmatrix}$ 
```

其中使用  $WY$  迭代来累积 Householder 变换的思想来源于以下算法:

#### (Matrix Computation 算法 5.1.2)

设有  $r \leq n$  个 Householder 变换  $H_1, \dots, H_r$ , 其中:

$$H_j = I_n - \beta_j v^{(j)} (v^{(j)})^H$$

$$v^{(j)} = [\underbrace{0, \dots, 0}_{j-1}, 1, v_{j+1}^{(j)}, \dots, v_n^{(j)}]^T$$

我们可以计算得到  $W, Y \in \mathbb{C}^{n \times r}$  满足  $H_1 \cdots H_r = I_n + WY^H$ :

$$\begin{aligned} Y &= v^{(1)} \\ W &= -\beta_1 v^{(1)} \\ \text{for } j &= 2 : r \\ z &= -\beta_j (I_n + WY^H)v^{(j)} = -\beta_j [v^{(j)} + W(Y^H v^{(j)})] \\ W &= [W, z] \\ Y &= [Y, v^{(j)}] \\ \text{end} \end{aligned}$$

$WY$  迭代累积  $n - 2$  个 Householder 变换相比直接累积  $n - 2$  个 Householder 变换并不会减少计算复杂度。但是可以增加 BLAS3 运算的比例，降低通讯开销，从而减少运行时间。

Matlab 代码为：

```
function [H, Q] = Hessenberg_Reduction_Householder_WY(A)
    % Hessenberg form using WY accumulation of Householder transformations
    % Input:
    % - A: Complex matrix of size n x n
    % Output:
    % - H: Upper Hessenberg form of A
    % - Q: Orthogonal matrix such that A = Q * H * Q'

    % Initialize matrices W and Y for WY accumulation, and vector b for storing betas
    [n, ~] = size(A);
    W = zeros(n-1, n-2); % Matrix to accumulate transformations for Q
    Y = zeros(n-1, n-2); % Matrix to store Householder vectors
    b = zeros(1, n-2); % Vector to store betas

    % Loop through each column (except the last two) for Householder reduction
    for k = 1:n-2
        % Step 1: Compute the Householder vector 'v' and scalar 'beta' for the current column
        [v, beta] = Complex_Householder(A(k+1:n, k));

        % Step 2: Apply the Householder transformation to zero out entries below the subdiagonal
        A(k+1:n, k:n) = A(k+1:n, k:n) - beta * v * (v' * A(k+1:n, k:n));
        A(1:n, k+1:n) = A(1:n, k+1:n) - (A(1:n, k+1:n) * v) * (beta * v)';

        % Step 3: Store the Householder vector in Y and the scalar beta in b
        Y(k:n-1, k) = v; % Store the Householder vector for later use
        b(k) = beta; % Store the scalar beta for the Householder transformation
    end

    % Loop to compute the W matrix, which accumulates the Householder transformations
    for k = 1:n-2
        % Step 4: Compute the W matrix for WY transformation
        if k == 1
            % For the first iteration, directly compute W
            W(1:n-1, 1) = -b(1) * Y(1:n-1, 1);
        else
            W(1:n-1, k) = W(1:n-1, 1:k-1) * (Y(k:n-1, 1:k-1)' * Y(k:n-1, k));
            W(k:n-1, k) = W(k:n-1, k) + Y(k:n-1, k);
            W(1:n-1, k) = -b(k) * W(1:n-1, k);
        end
    end

    % Step 5: Extract the upper Hessenberg matrix H from the modified matrix A
    H = A; % The matrix A is now in upper Hessenberg form after applying Householder

    % Step 6: Compute the orthogonal matrix Q as per WY transformation
    Q = eye(n, n);
    Q(2:n, 2:n) = Q(2:n, 2:n) + W * Y';
end
```

复数域上的 Householder 变换已于 Homework 02 Problem 02 给出：

```
function [v, beta] = Complex_Householder(x)
    % This function computes the Householder vector 'v' and scalar 'beta' for
    % a given complex vector 'x'. This transformation is used to create zeros
    % below the first element of 'x' by reflecting 'x' along a specific direction.

    n = length(x);
    x = x / norm(x, inf); % Normalize x by its infinity norm to avoid numerical issues
```

```

% Copy all elements of 'x' except the first into 'v'
v = zeros(n, 1);
v(2:n) = x(2:n);

% Compute sigma as the squared 2-norm of the elements of x starting from the second element
sigma = norm(x(2:n), 2)^2;

% Check if sigma is near zero, which would mean 'x' is already close to a scalar multiple of e_1
if sigma < 1e-10
    beta = 0; % If sigma is close to zero, set beta to zero (no transformation needed)
else
    % Determine gamma to account for the argument of complex number x(1)
    if abs(x(1)) < 1e-10
        gamma = 1; % If x(1) is close to zero, set gamma to 1
    else
        gamma = x(1) / abs(x(1)); % otherwise, set gamma to x(1) divided by its magnitude
    end

    % Compute alpha as the Euclidean norm of x, including x(1) and sigma
    alpha = sqrt(abs(x(1))^2 + sigma);

    % Compute the first element of 'v' to avoid numerical cancellation
    v(1) = -gamma * sigma / (abs(x(1)) + alpha);

    % Calculate 'beta', the scaling factor of the Householder transformation
    beta = 2 * abs(v(1))^2 / (abs(v(1))^2 + sigma);

    % Normalize the vector 'v' by v(1) to ensure that the first element is 1,
    % allowing for simplified storage and computation of the transformation
    v = v / v(1);
end
end

```

可视化正交性损失的函数:

```

function visualize_orthogonality_loss(Q, titlestr)
    % Visualizes the componentwise loss of orthogonality |Q^H Q - I_n|
    loss = Q' * Q - eye(size(Q, 2)); % Compute the loss
    figure; % Create a new figure window
    imagesc(log10(abs(loss))); % Display the absolute value of the loss
    colorbar; % Add colorbar to indicate scale
    title(titlestr);
    xlabel('Column Index');
    ylabel('Row Index');
    axis square; % Make the axes square for better visualization
end

```

函数调用:

```

rng(51);
n = 100;
A = rand(n, n) + 1i * rand(n, n);

% Apply the Hessenberg reduction
[H, Q] = Hessenberg_Reduction_Householder_WY(A);

% Display the Frobenius norm of Q' * A * Q - H
disp("Frobenius norm of Q' * A * Q - H:");
disp(norm(Q' * A * Q - H, 'fro')); % Compute Frobenius norm for forward error

% Display the Frobenius norm of Q' * Q - In
disp("Frobenius norm of Q' * Q - In:");
disp(norm(Q' * Q - eye(n), 'fro'));

% Visualize the loss of orthogonality for Q
visualize_orthogonality_loss(Q, 'Log10 Loss of Orthogonality (Householder Hessenberg Reduction)');

% Perturb the matrix A slightly and repeat the check
% A_perturbed is a slightly perturbed version of A, with a small random perturbation added to each element.
A_perturbed = A + 1e-5 * (randn(n, n) + 1i * randn(n, n)); % Perturb A by a small amount

% Apply the Hessenberg reduction to the perturbed matrix
[H_perturbed, Q_perturbed] = Hessenberg_Reduction_Householder_WY(A_perturbed);

% Check the accuracy for the perturbed case
% Compute and display the forward error between the perturbed and original Hessenberg matrices.
% This tells us how much the Hessenberg matrix changes due to the perturbation in A.

```

```

disp("Frobenius norm of H_perturbed - H: (Forward Error)");
disp(norm(H_perturbed - H, 'fro') / norm(A_perturbed - A, 'fro')); % Forward error for Hessenberg matrix

% Compute and display the forward error for the orthogonal matrix Q
% This tells us how much the orthogonal matrix Q changes due to the perturbation in A.
disp("Frobenius norm of Q_perturbed - Q: (Forward Error)");
disp(norm(Q_perturbed - Q, 'fro') / norm(A_perturbed - A, 'fro')); % Forward error for Q

% Compute and display the backward error
% This checks how well the perturbed matrix Q_perturbed and H_perturbed approximate the original matrix A.
% It is a backward error computation that checks the closeness of Q_perturbed * H_perturbed * Q_perturbed' to A.
disp("Frobenius norm of Q' * A * Q - H: (Backward Error)");
disp(norm(Q_perturbed * H_perturbed * Q_perturbed' - A, 'fro') / norm(A_perturbed - A, 'fro')); % Backward error for Hessenberg matrix approximation

```

运行结果:

```

Frobenius norm of Q' * A * Q - H:
1.0942e-13

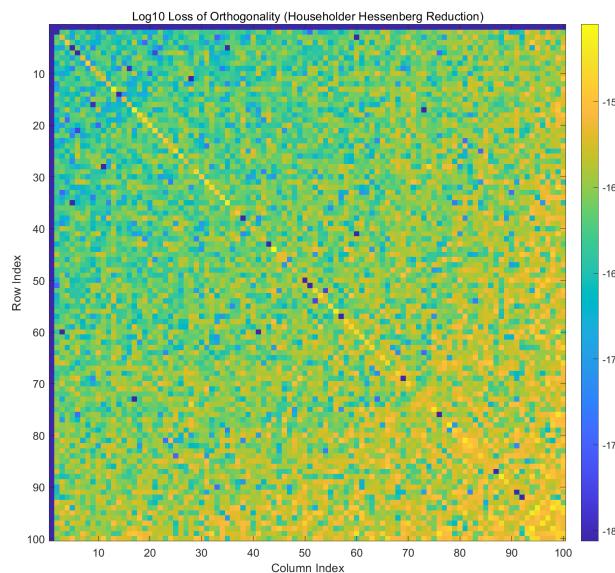
Frobenius norm of Q' * Q - In:
1.3307e-14

Frobenius norm of H_perturbed - H: (Forward Error)
1.0585e+03

Frobenius norm of Q_perturbed - Q: (Forward Error)
248.1138

Frobenius norm of Q' * A * Q - H: (Backward Error)
1.0000

```



## Part (2)

基于 Gram-Schmidt 正交化实现 Arnoldi 过程的函数已于 Homework 6 Problem 4 给出:  
(我们只需将  $r$  设为  $n$  即可对  $A$  进行上 Hessenberg 化)

```

function [Q, H] = Gram_Schmidt_Arnoldi(A, b, r, tolerance, modified, reorthogonalized)
    % Gram_Schmidt_Arnoldi computes an orthonormal basis Q and a Hessenberg matrix H
    % using the Arnoldi process with either Classical or Modified Gram-Schmidt.
    %
    % Inputs:
    %   A - Square matrix of size n x n.
    %   b - Initial vector of size n x 1.
    %   r - Desired rank of the output.
    %   tolerance - Threshold for detecting linear dependence (default: 1e-10).
    %   modified - Flag for using Modified Gram-Schmidt (default: false).
    %   reorthogonalized - Flag for reorthogonalization (default: false).
    %
    % Outputs:

```

```

% Q - Orthonormal basis of size n x r.
% H - Upper Hessenberg matrix of size r x r.

% Validate inputs and set default values if not provided
if nargin < 6
    reorthogonalized = false; % Default to not reorthogonalizing
end
if nargin < 5
    modified = false; % Default to Classical Gram-Schmidt
end
if nargin < 4
    tolerance = 1e-10; % Default tolerance for linear dependence
end

% Get the size of the matrix A
n = size(A, 1);

% Check if the desired rank is valid
if r < 1 || r > n
    error("r should be an integer in [1, n]");
end

% Initialize matrices Q and H
Q = zeros(n, n);
Q(:, 1) = b / norm(b); % Normalize the initial vector b
H = zeros(r, r); % Initialize H to zeros
delta = zeros(n, 1); % Temporary variable for inner products

% Set max iterations based on reorthogonalization flag
if reorthogonalized
    max_iter = 2; % More iterations for reorthogonalization
else
    max_iter = 1; % Standard iteration count
end

% Loop to build the orthonormal basis up to r-1 or n-1
for k = 1:r-1
    % Apply the matrix A to the last basis vector
    Q(:, k + 1) = A * Q(:, k);

    if modified
        % Modified Gram-Schmidt process
        for iter = 1:max_iter
            for i = 1:k
                % Compute inner product
                delta(i) = Q(:, i)' * Q(:, k + 1);
                % Update H matrix
                H(i, k) = H(i, k) + delta(i);
                % Orthogonalize the k+1 vector
                Q(:, k + 1) = Q(:, k + 1) - delta(i) * Q(:, i);
            end
        end
    else
        % Classical Gram-Schmidt process
        for iter = 1:max_iter
            % Compute inner products for all previous basis vectors
            delta(1:k) = Q(:, 1:k)' * Q(:, k + 1);
            % Update H matrix
            H(1:k, k) = H(1:k, k) + delta(1:k);
            % Orthogonalize the k+1 vector
            Q(:, k + 1) = Q(:, k + 1) - Q(:, 1:k) * delta(1:k);
        end
    end

    % Compute the norm for the current basis vector
    H(k + 1, k) = norm(Q(:, k + 1));

    % Check for linear dependence by comparing the norm with the tolerance
    if H(k + 1, k) < tolerance
        fprintf("The rank %d is lesser than %d\n", k, r);
        r = k; % Update the rank if linear dependence is detected
        break; % Exit the loop early
    else
        % Normalize the current basis vector
        Q(:, k + 1) = Q(:, k + 1) / H(k + 1, k);
    end
end

% Fill the last column of H

```

```

H(1:r, r) = Q(:, 1:r)' * (A * Q(:, r));

% Trim Q and H to the computed effective rank
Q = Q(:, 1:r);
H = H(1:r, 1:r);
end

```

函数调用:

```

rng(51);
n = 100;
A = rand(n, n) + 1i * rand(n, n);
b = rand(n, 1) + 1i * rand(n, 1);

% Apply the Hessenberg reduction
% [Q, H] = Gram_Schmidt_Arnoldi(A, b, r, tolerance, modified, reorthogonalized)
[Q, H] = Gram_Schmidt_Arnoldi(A, b, n, 1e-10, true, true);

% Display the Frobenius norm of Q' * A * Q - H
disp("Frobenius norm of Q' * A * Q - H:");
disp(norm(Q' * A * Q - H, 'fro')); % Compute Frobenius norm for forward error

% Display the Frobenius norm of Q' * Q - In
disp("Frobenius norm of Q' * Q - In:");
disp(norm(Q' * Q - eye(n), 'fro'));

% Visualize the loss of orthogonality for Q
visualize_orthogonality_loss(Q, 'Log10 Loss of Orthogonality (Householder Hessenberg Reduction)');

% Perturb the matrix A slightly and repeat the check
% A_perturbed is a slightly perturbed version of A, with a small random perturbation added to each element.
A_perturbed = A + 1e-5 * (randn(n, n) + 1i * randn(n, n)); % Perturb A by a small amount

% Apply the Hessenberg reduction to the perturbed matrix
% [Q, H] = Gram_Schmidt_Arnoldi(A, b, r, tolerance, modified, reorthogonalized)
[Q_perturbed, H_perturbed] = Gram_Schmidt_Arnoldi(A_perturbed, b, n, 1e-10, true, true);

% Check the accuracy for the perturbed case
% Compute and display the forward error between the perturbed and original Hessenberg matrices.
% This tells us how much the Hessenberg matrix changes due to the perturbation in A.
disp("Frobenius norm of H_perturbed - H: (Forward Error)");
disp(norm(H_perturbed - H, 'fro') / norm(A_perturbed - A, 'fro')); % Forward error for Hessenberg matrix

% Compute and display the forward error for the orthogonal matrix Q
% This tells us how much the orthogonal matrix Q changes due to the perturbation in A.
disp("Frobenius norm of Q_perturbed - Q: (Forward Error)");
disp(norm(Q_perturbed - Q, 'fro') / norm(A_perturbed - A, 'fro')); % Forward error for Q

% Compute and display the backward error
% This checks how well the perturbed matrix Q_perturbed and H_perturbed approximate the original matrix A.
% It is a backward error computation that checks the closeness of Q_perturbed * H_perturbed * Q_perturbed' to A.
disp("Frobenius norm of Q' * A * Q - H: (Backward Error)");
disp(norm(Q_perturbed * H_perturbed * Q_perturbed' - A, 'fro') / norm(A_perturbed - A, 'fro')); % Backward error for Hessenberg matrix approximation

```

运行结果:

```

Frobenius norm of Q' * A * Q - H:
5.3483e-14

Frobenius norm of Q' * Q - In:
4.5518e-15

Frobenius norm of H_perturbed - H: (Forward Error)
34.0533

Frobenius norm of Q_perturbed - Q: (Forward Error)
8.0111

Frobenius norm of Q' * A * Q - H: (Backward Error)
1.0000

```

