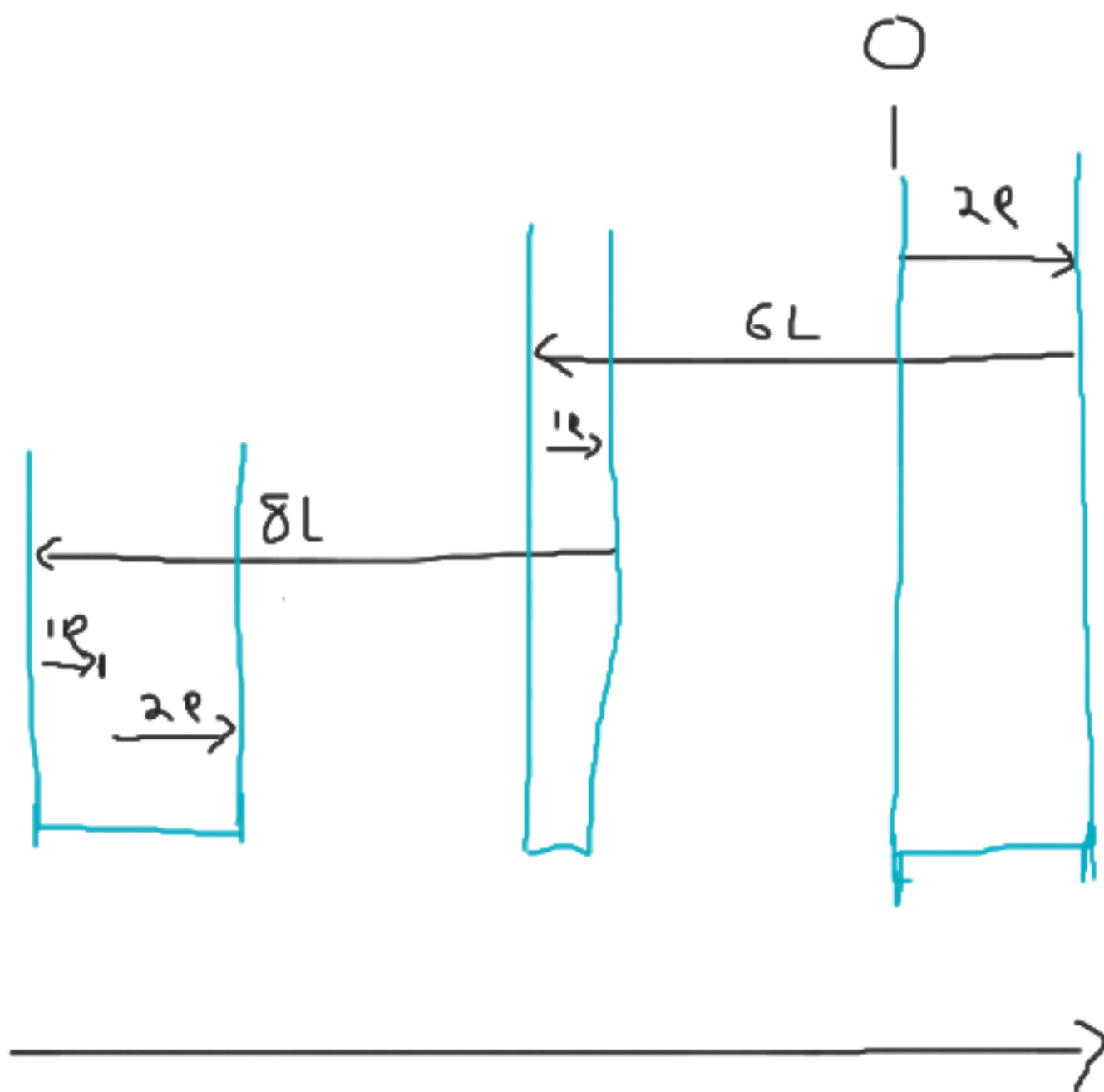


Self-Reflection Form: Until 9:30

6. Painting the Fence



$(0, 2)$
 $(2, -4)$
 $(-4, -3)$
 $(-3, -11)$
 $(-11, -10)$
 $(-10, -8)$

Runtime $O(N \log N)$

(S) (E)
 (-11) (-10)
 $(-11, -10)$
 $(-11, 3)$ (-3)
 $(-10, -8)$
 $(-4, -4)$
 $(-4, 2)$
 $(0, 2)$

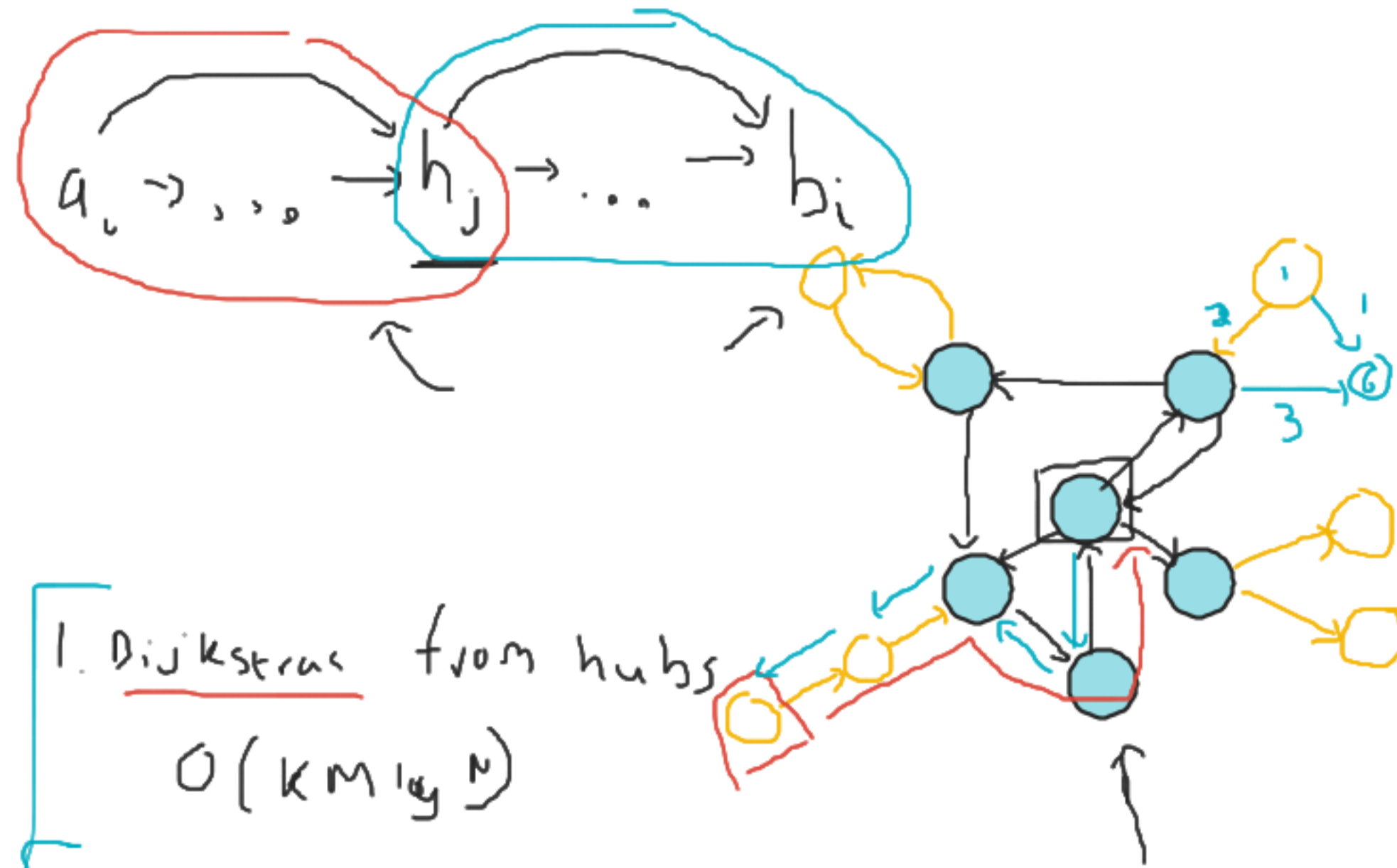
counts

ans

(S) 1
 (-11) 2
 (E) 1
 (-10)

(S) 2 1
 (-10) 1 3
 (E) 1 3
 (-8)

3. Vacation Planning



1. Dijkstra's from hubs

$$O(kM \log N)$$

2. Reverse then (1)

$$O(kM \log N)$$

Naive
 $O(QM \log N)$

Floyd-Warshall

$$O(N^3 + QK)$$

Dijkstra's

$$O(kM \log N + QK)$$

$$1 \rightarrow 6$$

$$\begin{matrix} h \\ 2 \\ 3 \end{matrix} \text{dist}[1][2] \quad \text{dist}[2][1]$$

4. A Coin Game

$\overset{n}{\downarrow} \overset{n}{\downarrow}$
 $DP[i][j] = \text{max score at position } i \text{ picking up } j \text{ coins}$

if $i + j > n$

$DP[i][j] = \text{sum of all coins after } i,$

else

$$DP[i][j] = \max \left(DP[i+j][\min(j*2, N)], DP[i][j-1] \right)$$

sig = 0
 for $i: N-1 \rightarrow 0$

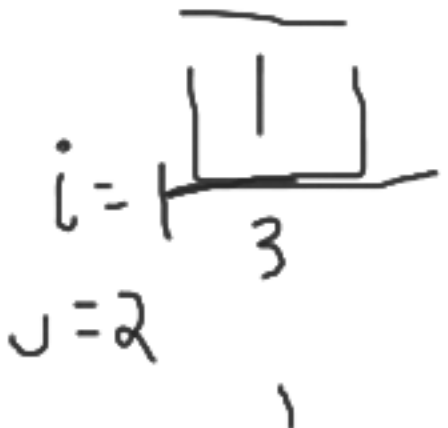
sig += $d[i+1]$

for $j: 1 \rightarrow N$

$\max(\text{sig} - dp[i+j][\min(j*2, N)], dp[i][j-1])$

$\leq j$

$O(N)$
 prefix



Break 1:	10:30 AM - 10:45 AM
Lunch:	12:15 PM - 1:30 PM
Break 2:	3:00 PM - 3:15 PM
End:	4:45 PM

$$f(i) = f(i-1) + f(i-2)$$

d_p init - 1

$$f(\phi) = 1$$
$$f(a) = 1$$
$$\text{def } f(i)$$

→ if $i \leq 2$

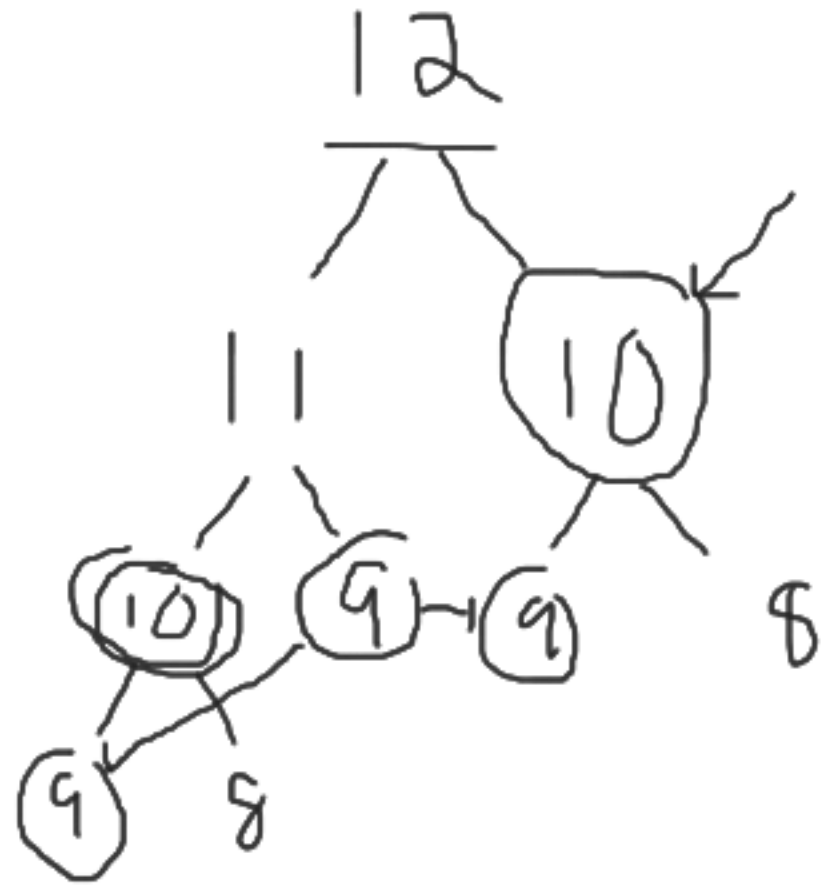
return 1

```
dp[i] = -1
return dp[i]
```

$e|_x$

$$\text{dp}[i] = \text{return } f(i-1) + f(i-2)$$

return dp[i]

 $O(N)$

knapsack

11 lbs

weights

3 lbs ∞

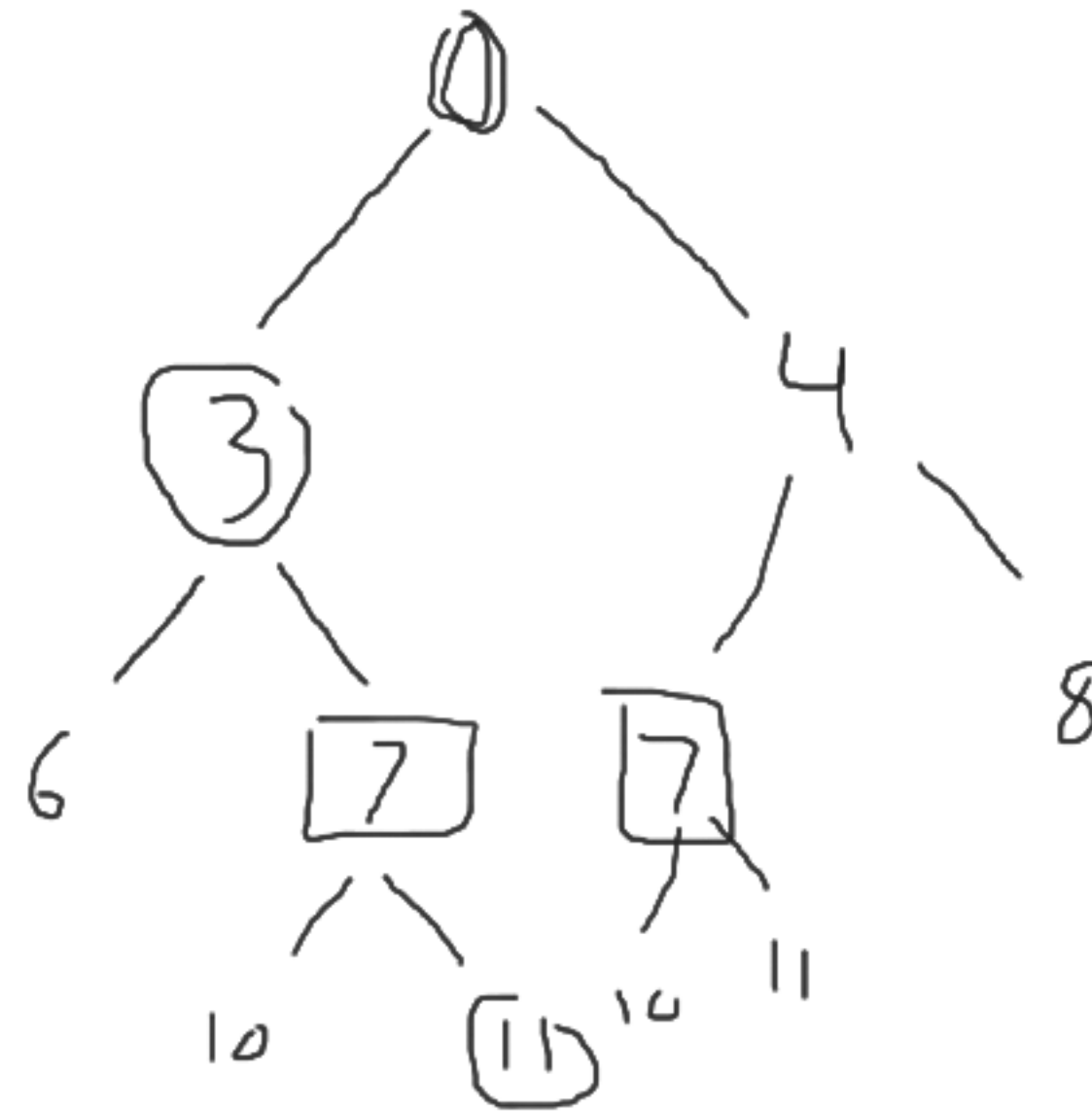
4 lbs ∞

$dp[0] = true$ ans
for $i: 0 \rightarrow 11$

if $dp[i]$:

$dp[i+3] = true$

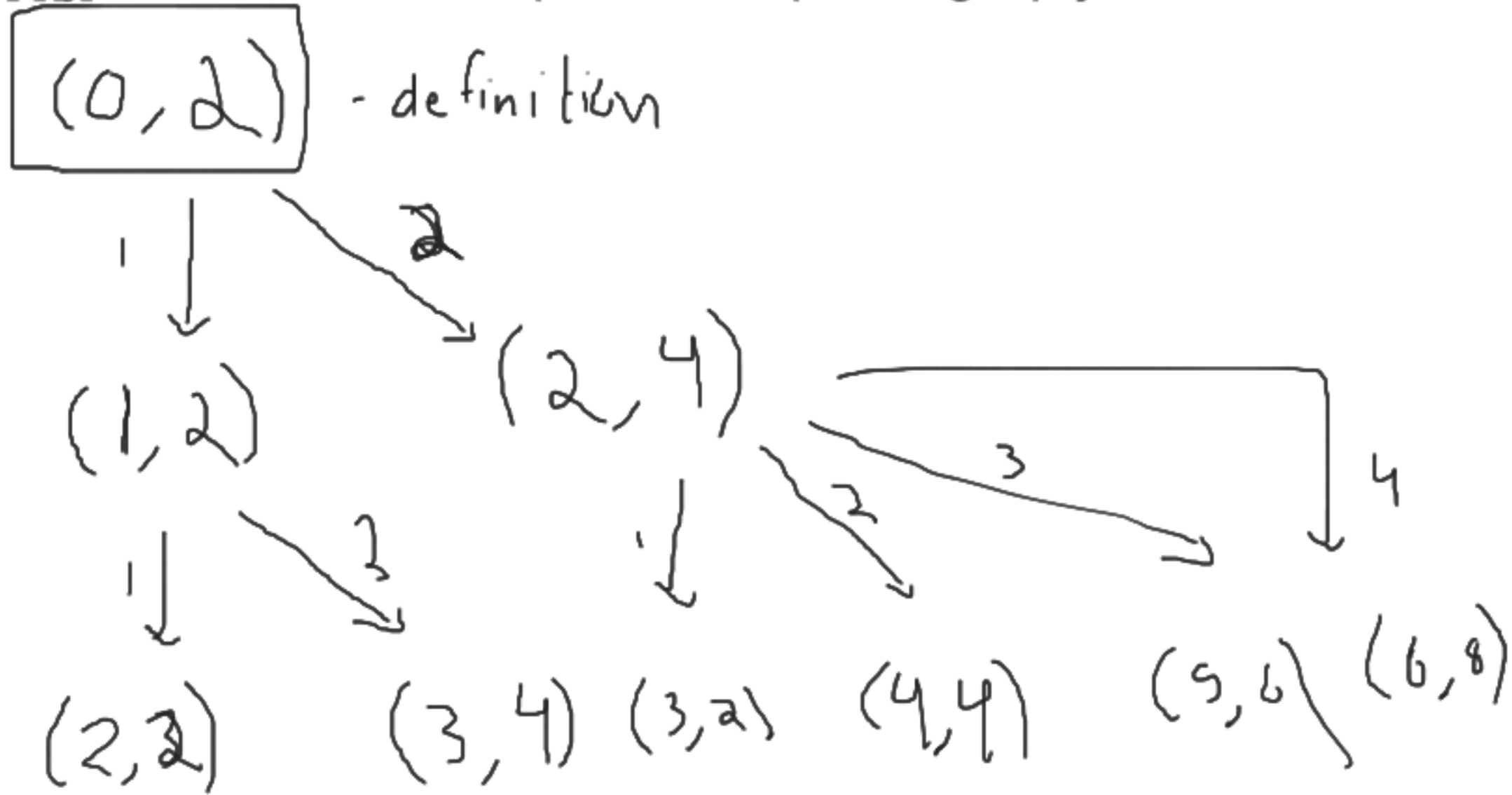
$dp[i+4] = true$



$dp[i] = \text{can we do weight } i / \text{true or false}$

$dp[i] \Rightarrow dp[i+3]$
 $\Rightarrow dp[i+4]$

$DP[i][j]$ = max score at position i picking up j coins



Tile Exchanging

DP STATE

$$DP[i][j] = \text{Area } i \text{ Tiles } 1 \dots j \rightarrow \text{min cost}$$

BASE CASE

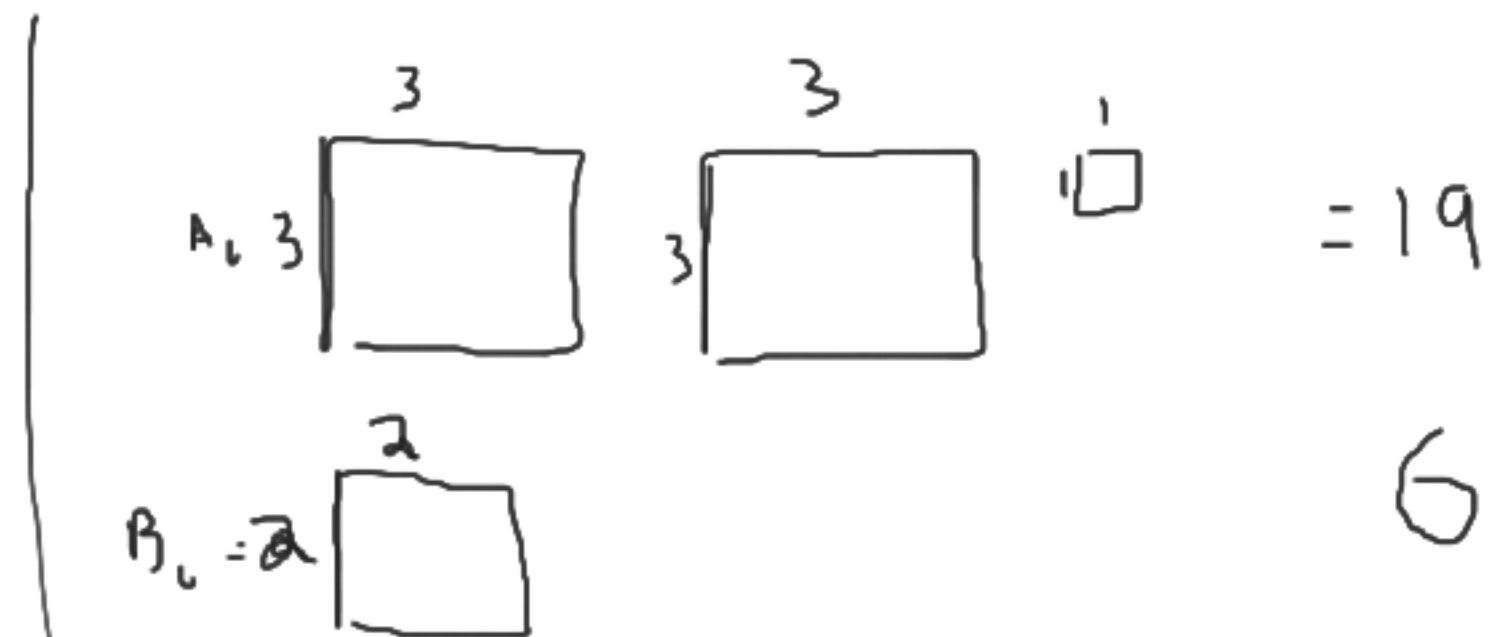
$$DP[0][0] = 0$$

TRANSITION

$$DP[i][j] = \min_k \left[\underbrace{dp[i - k^2][j - 1]}_{\text{with } j-1 \text{ tiles}} \right] + (A_j - k)^2$$

RUNTIME

$$O\left(\underbrace{M \cdot N \cdot \sqrt{M}}_{\text{dp state}}\right) = 10^4 \cdot 10^1 \cdot 10^2 = 10^7$$



$$|A_i - B_i|, |A_i - B_i|^2 = (A_i - B_i)^2$$

$$N \leq 10$$

$$M \leq 10,000$$

$$1 \xrightarrow{1} 2 \xrightarrow{1} 3 \quad 1 \xrightarrow{4} 3$$

↖ ↗

↖ ↗

↓ ↘ ↗

✓	✓		
✓	✓		
✓	✓		
✓	✓		
✓	✓		
✓	✓		



$$\underline{i = 6}, \underline{j = 3}$$

Case 1 1 

$$dp[s][2]$$

Case 2 2 

$$dp[2][2] + (2-1)^2$$

Case 3 3 

5 3

1 2



1 3

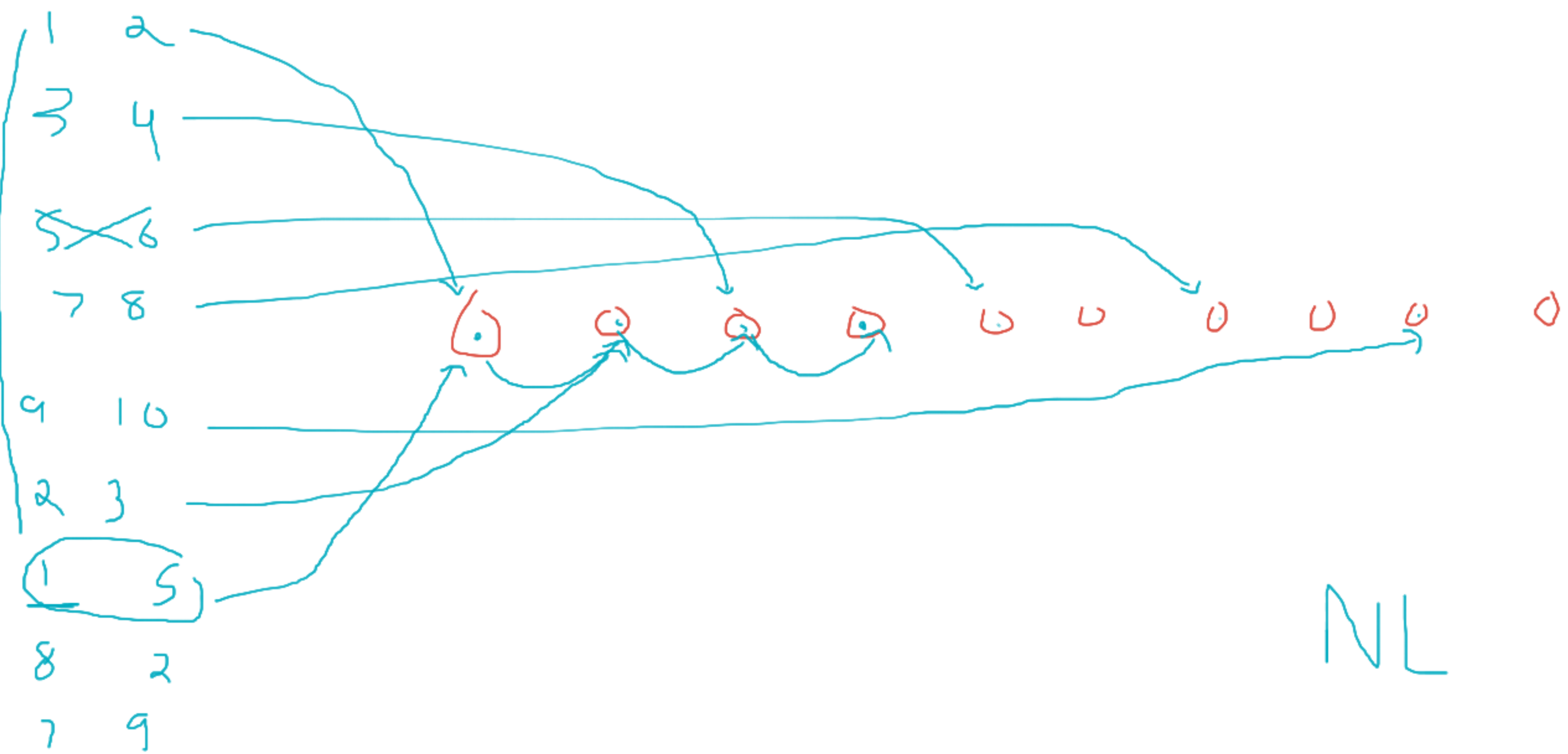
1 2

1 3

1 2

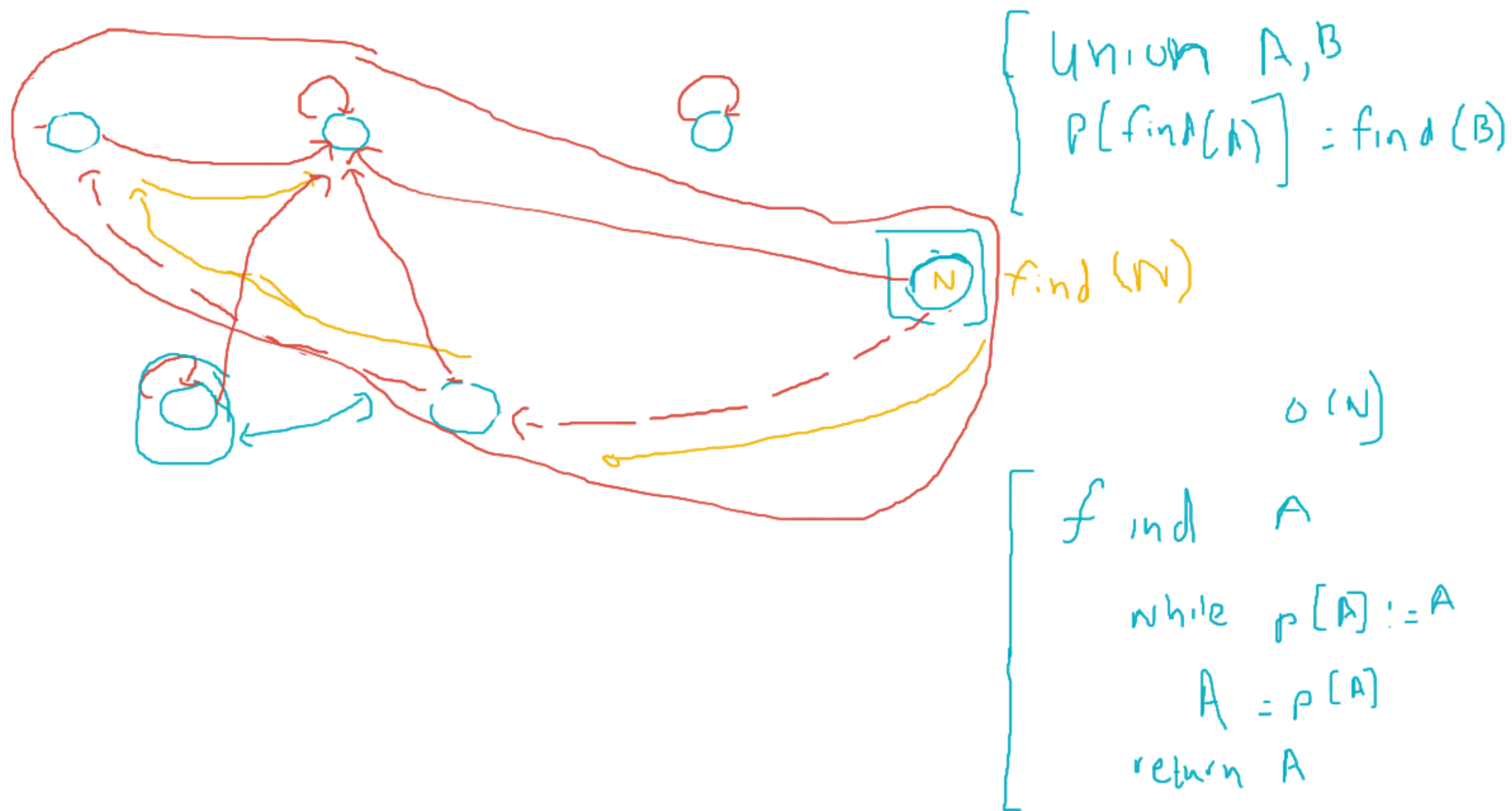


1.  If the drawer A_i is empty, he stores the item i in that drawer
2.  If the drawer B_i is empty, he stores the item i in that drawer
3. Try to move the item from A_i to its other drawer; if that one's filled too, try moving that item to its other drawer, and so on until you either succeed or get back to a previously seen drawer. In case of success, store the item i in the drawer A_i . In case of failure, continue to next rule.
4. Try moving the item from B_i to its other drawer; if that one's filled too, try moving that item to its other drawer, and so on until you either succeed or get back to a previously seen drawer. In case of success, store the item i in the drawer B_i . In case of failure, continue to next rule.
5. Give up and throw away the item i . For given pairs of drawers for each item, determine which items will be stored and which will be thrown away.

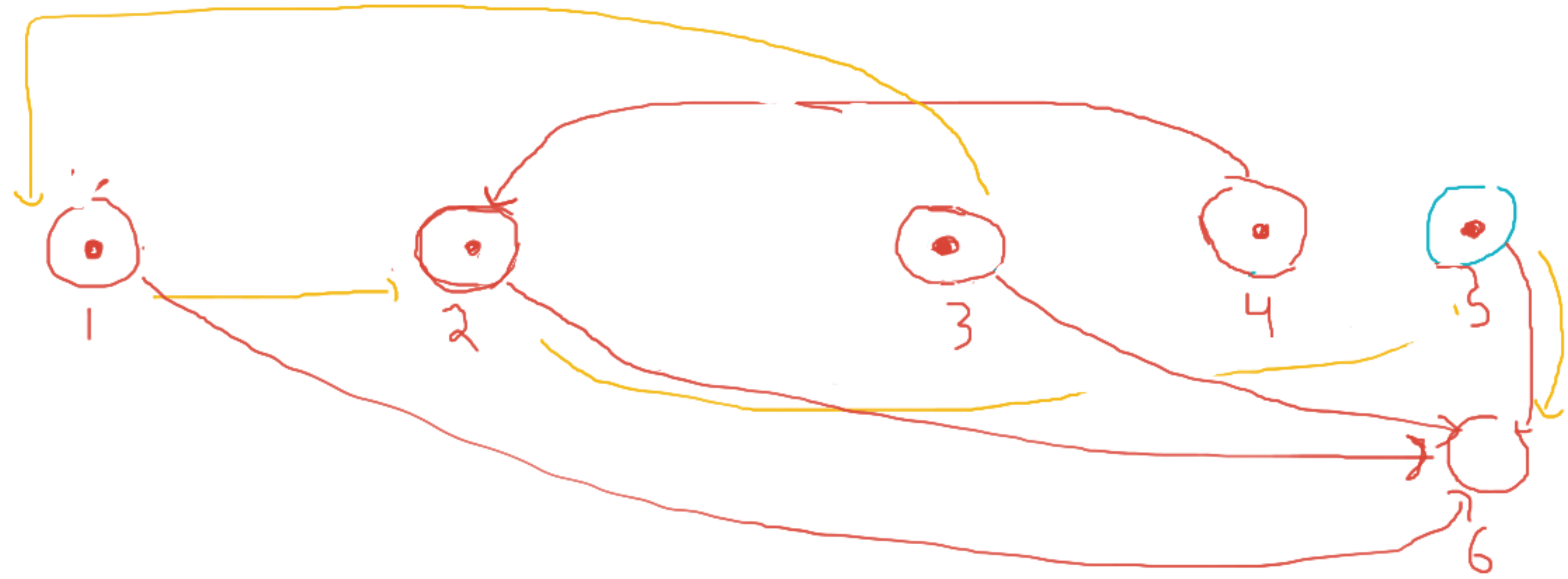


NL

$O(N)$



1,2
1,3
4,2
3,5
...



$O(NL)$

$$A = \begin{bmatrix} \begin{matrix} \times \\ 2 \\ 3 \\ 4 \\ 5 \end{matrix} \end{bmatrix}$$

$$\begin{matrix} 2 \\ 1 \\ 3 \\ 4 \\ 5 \end{matrix}$$

$$A =$$

$$\begin{bmatrix} 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

$$\boxed{1 \ 3 \ 4 \ 5 \ 2}$$

$$\rightarrow 1 \ 3 \ 4 \ 5 \ 2$$

$$\boxed{2 \ [1 \ 3 \ 4] \ 5}$$

[1 3 4] s 2

$$\begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

shuffle
+ left

$$= \begin{matrix} & A \\ \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix} \end{matrix} \begin{bmatrix} 1 \\ 2 \\ 3 \end{bmatrix} = \begin{bmatrix} 1 \cdot 0 + 2 \cdot 1 + 3 \cdot 0 \\ 1 \cdot 1 + 2 \cdot 0 + 3 \cdot 0 \\ 1 \cdot 0 + 2 \cdot 0 + 3 \cdot 1 \end{bmatrix} = \begin{bmatrix} 2 \\ 1 \\ 3 \end{bmatrix}$$

$$= \begin{bmatrix} 1 \\ 2 \\ 3 \end{bmatrix}$$

$$A^N \cdot \begin{bmatrix} 1 \\ 2 \\ 3 \end{bmatrix}$$

$$2^7 \quad \swarrow \quad 0b \overset{4}{1} \overset{2}{1} \overset{1}{1}$$

$$4 + 2 + 1 = 7$$

$$2^7 = 2^4 \cdot 2^2 \cdot 2^1$$

$$2^2 = 2^1 \cdot 2^1$$

$$2^4 = 2^2 \cdot 2^2$$

$$2^5 \quad \nearrow \quad 0b \underline{1} \underline{0} \underline{1}$$

