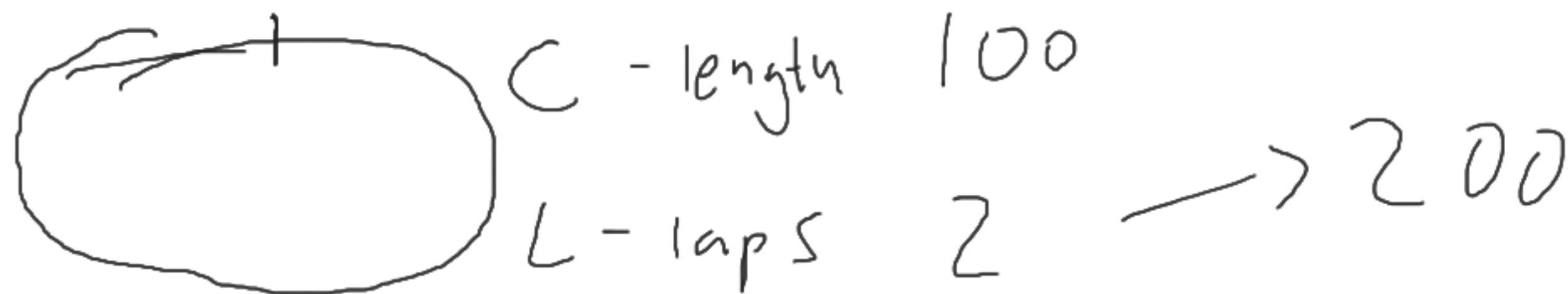


N cows



20, 100, 70, 1

$$200/100 = 2 \text{ units of time}$$

Sample output: 4

100: 20, 1

70: 20, 1

$l(i)$ = Number of laps done by cow i
during the entire race $\in \mathbb{R}$

$$l(i) = \frac{\text{Speed}_i \cdot \text{time}_T}{C}$$

$\text{time}_T = \frac{C \cdot L}{\max(\text{speed})}$

1	20	70	100
0.02	0.4	1.4	2

cows j, k How many times does k overtake j

$$\text{floor}(l(k) - l(j))$$

25

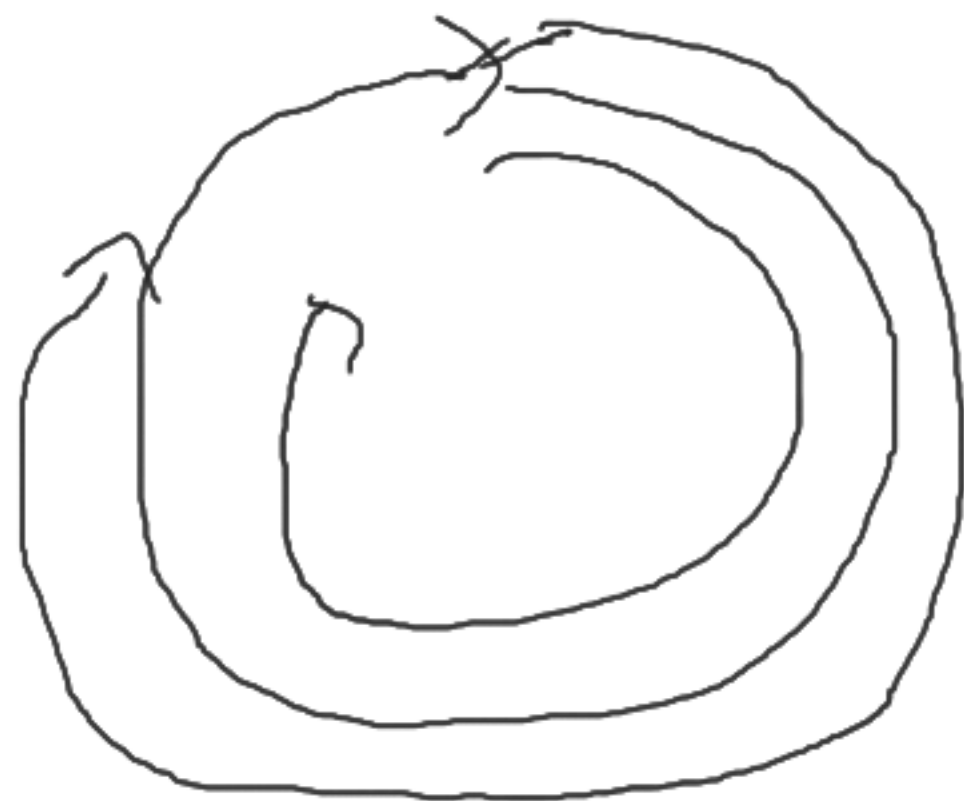
50

100

1

2

4



1.8

0.8

$$\sum [l(b) - l(a)] \text{ for all } b > a$$

$$\sum l(b) - l(a) \text{ for all } b > a$$

$$[1] \quad [1] \quad [2] \quad [4] \rightarrow [1 - (1)] + [2 \cdot 2 - (1+1)] + [4 \cdot 3 - (2+1+1)]$$

$$[1 - (1)] + [(2 - 1) + (2 - 1)] + [(4 - 2) + (4 - 1)]$$

Naive: $O(N^2)$

$$\sum l(b) - l(a) \text{ for } b > a$$

$$l(1)=1 \quad l(2)=1 \quad l(3)=2 \quad l(4)=4$$

$$[l(2) - l(1)] + [l(3) - l(2)] + [l(3) - l(1)]$$

$$+ [l(4) - l(3)] + [l(4) - l(2)] + [l(4) - l(1)]$$

$$= \overset{\delta_0}{[l(2) - \cancel{l(1)}]} + \overset{\delta_1}{[2 \cdot l(3) - (\cancel{l(2)} + \cancel{l(1)})]} + \overset{\delta_2}{[3 \cdot l(4) - (\cancel{l(3)} + \cancel{l(2)} + \cancel{l(1)})]}$$

$\delta_i = l(i+1) + \delta_{i-1}$

$$\ell = [1.2 \quad 1.3 \quad 2.2 \quad 4.4]$$

$$[\ell(2) - \ell(1)] = 0$$

$$[\ell(3) - \ell(2)] + [\ell(3) - \ell(1)] = 0 + 1 = 1$$

$$[\ell(4) - \ell(3)] + \dots = 2 + 3 + 3 = \frac{8}{9}$$

$$l = [1.2 \quad 1.3 \quad 2.2 \quad 4.4] \quad \left[\sum l(b) - l(a) \right]$$

$$[l(2)] - [l(1)] = 0$$

$$([l(3)] - [l(2)]) + ([l(3)] - [l(1)]) = 1 + 1 = 2$$

$$([l(4)] - [l(3)]) + \dots = 2 + 3 + 3 = \frac{8}{10}$$

1.3, 2.2 \rightarrow problem: $L\{l(b) - l(a)\} \neq$

$$\{x\} = x - Lx$$

$$L\{l(b)\} - L\{l(a)\}$$

$$\{l(a)\} > \{l(b)\}$$

$$\{1.3\} = 0.3$$

Binary-Index Tree / Fenwick Tree / BIT

$a = [\underbrace{0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0}_{\text{array}}]$

Array $O(1)$ | BIT $O(\log N)$ $\text{update}(i, val)$ — change the value at index i

$O(N)$ | $O(\log N)$ $\text{query}(i)$ = returns the sum of $\sum_{j=0}^i a[j]$

$p = [1.2 \quad 1.3 \quad 2.2 \quad 4.4]$

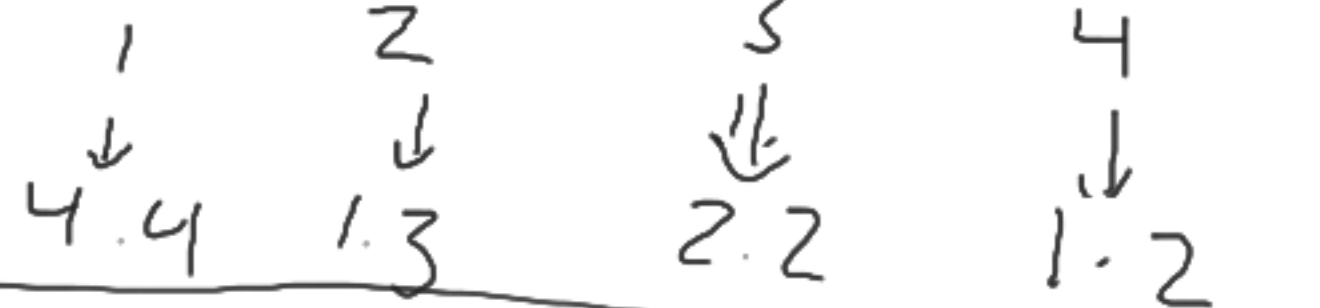
index

1: update($f(1.2)$, 1)

$O(N \log N)$

Runtime

BIT



$$2: \lfloor l(2) \rfloor - \lfloor l(1) \rfloor = 0$$

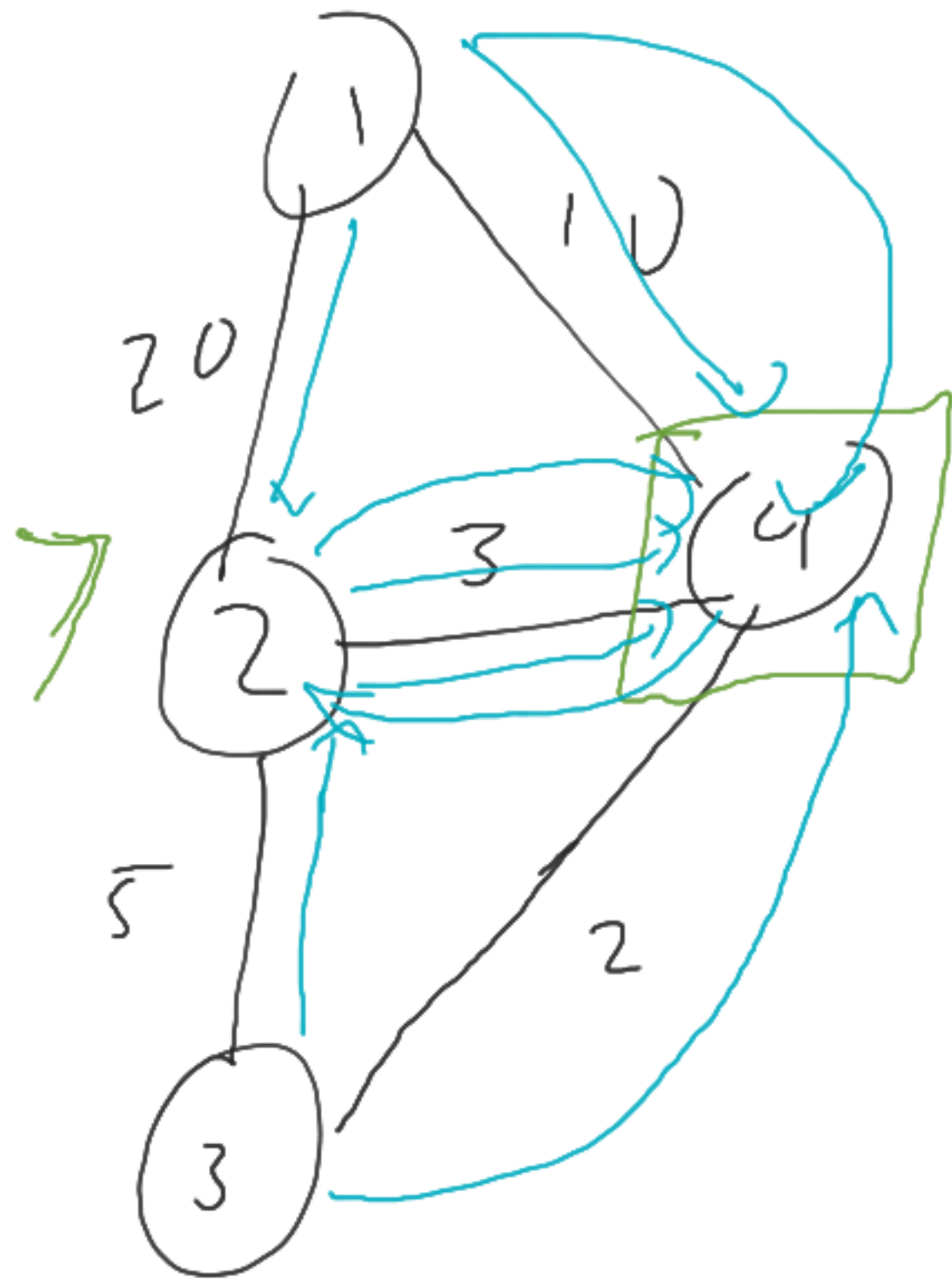
errors: sum of all elements before 1.3 = 0, update($f(1.3)$, 1)

$$3: (\lfloor l(3) \rfloor - \lfloor l(2) \rfloor) + (\lfloor l(3) \rfloor - \lfloor l(1) \rfloor) = 1 + 1 = 2$$

errors: sum of all elements before 2.2 = 1, update($f(2.2)$, 1)

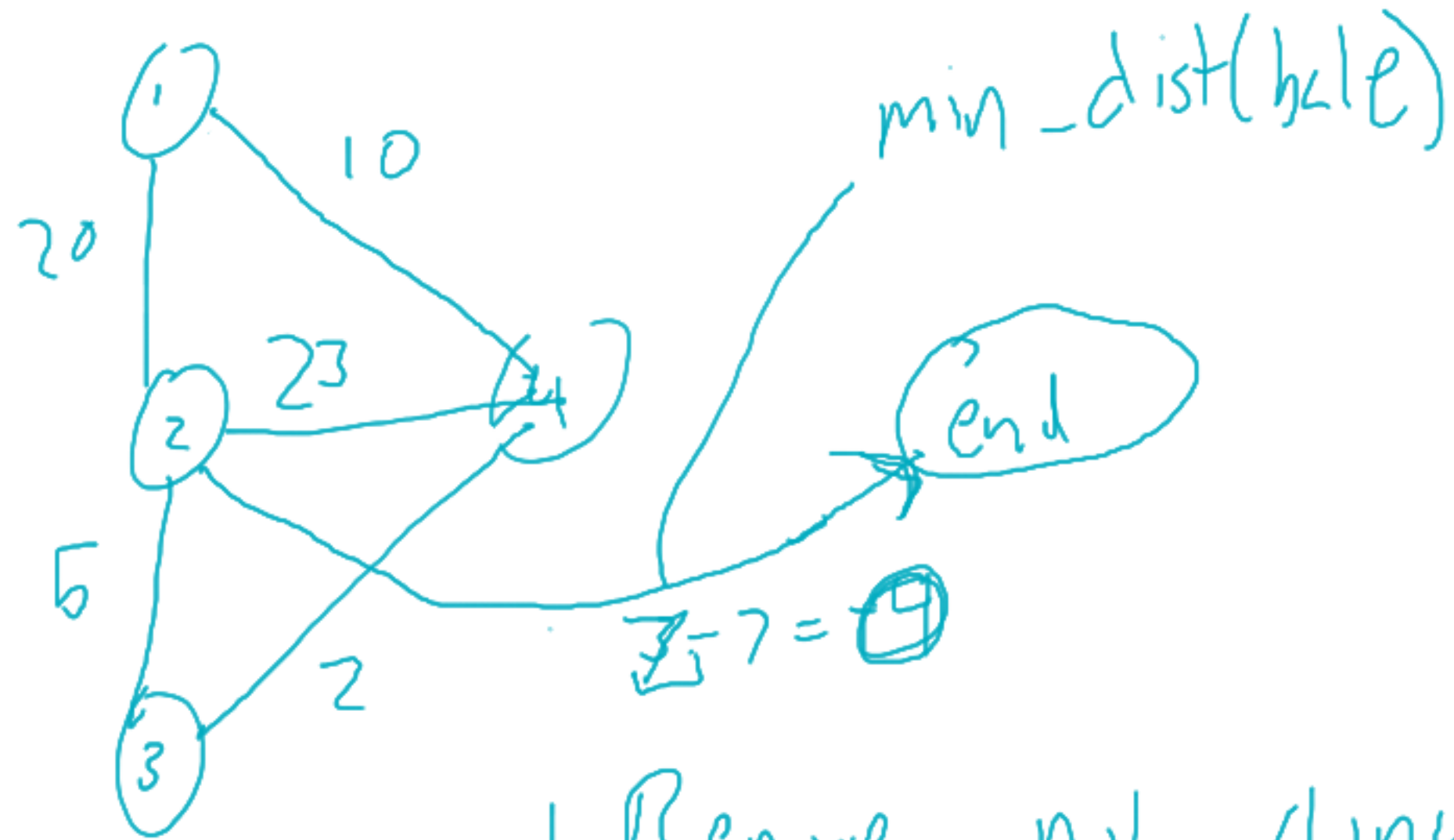
$$4: \lfloor l(4) \rfloor - \dots = 8 \quad \text{errors} = 0$$

$$\sum \lfloor l(4) \rfloor - \lfloor l(1) \rfloor \text{ errors} = 9$$



	Original 10	w/ thresholds 23 16	✓
3	2	8	✓

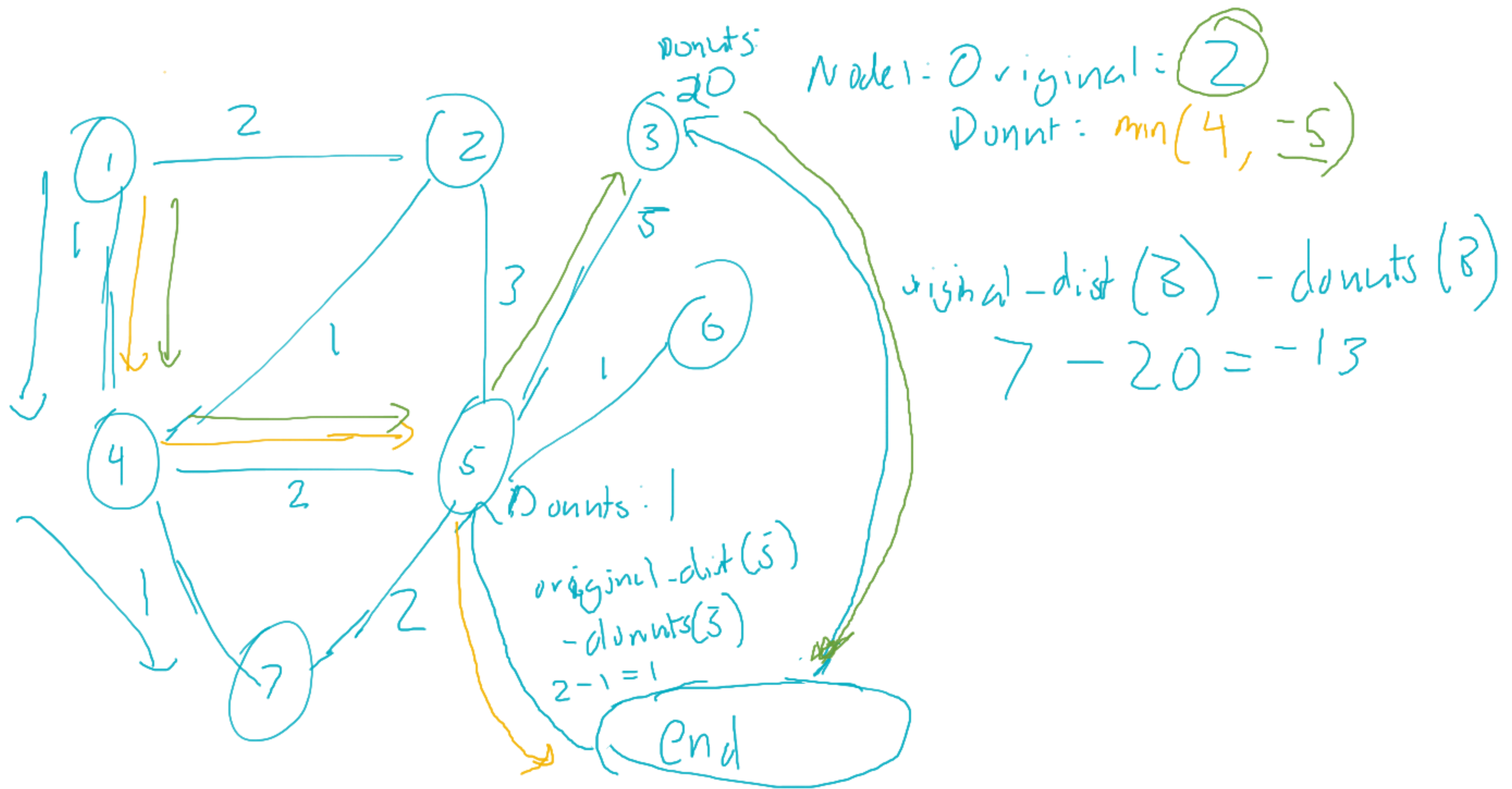
↑
 Dijkstra's from
 node N
 $\mathcal{O}(M \log N)$



Goals:

1. Force the cow to actually get donuts
2. Reduce the path weight by donut yumminess

1. Remove not donut edges
+ add an edge to every donut



```
public class BIT {

    int[] tree;
    int N;

    public BIT(int N) {
        this.N = N;
        tree = new int[N+1];
    }

    public BIT(int N, int[] d){
        this.N = N;
        tree = new int[N+1];
        for (int i = 1; i < d.length; i++) {
            update(i,d[i]);
        }
    }

    public int query(int K) {
        int sum = 0;
        for (int i = K; i > 0; i -= (i & -i)) {
            sum += tree[i];
        }
        return sum;
    }

    public void update(int K, int val) {
        for (int i = K; i <= N; i += (i & -i)) {
            tree[i] += val;
        }
    }
}
```

T 30

$1 \rightarrow \frac{1}{2} \rightarrow \frac{1}{3} \rightarrow \frac{1}{4}$

D 10

1000

T 10

T 20

T 50

dist 10

15

25



time 10

20

50

dist = 0 → 10 → 12.5

time = 0 → 10 → 15

Speed-denom = 1, 0 → 2.0 → 3

$O(N^2)$

$O(N \log N)$

²⁰
T 30 ~~D 10~~
¹⁰
T 20 D 55⁹⁰
~~T 15~~⁵ D 20²⁰

~~[T 15]~~
~~[T 20]~~
[T 30]

~~[D 10]~~
~~[D 20]~~
[D 55]

$dp[i] = \text{min cost to protect cows } 1 \dots i$
 $= f$

$f(i) =$
 $\text{ans} = \infty$
for l in $\text{range}(1, M)$: $\swarrow O(N^2 M)$
 $j := \text{first cow covered if we end an umbrella of length } l \text{ at } i$

$\text{ans} = \min(\text{ans}, f(j-1) + C(l))$

$dp[i] = \text{min cost of covering cows } 1 \dots i$

$$dp[i] = dp[i-1] + c[i]$$

for every $k < i$

$$dp[i] = \min(dp[i], dp[k] + c'[X_i - X_{k+1}])$$

$O(N^2)$

min cost of single umbrella covering cows $k+1 \dots i$

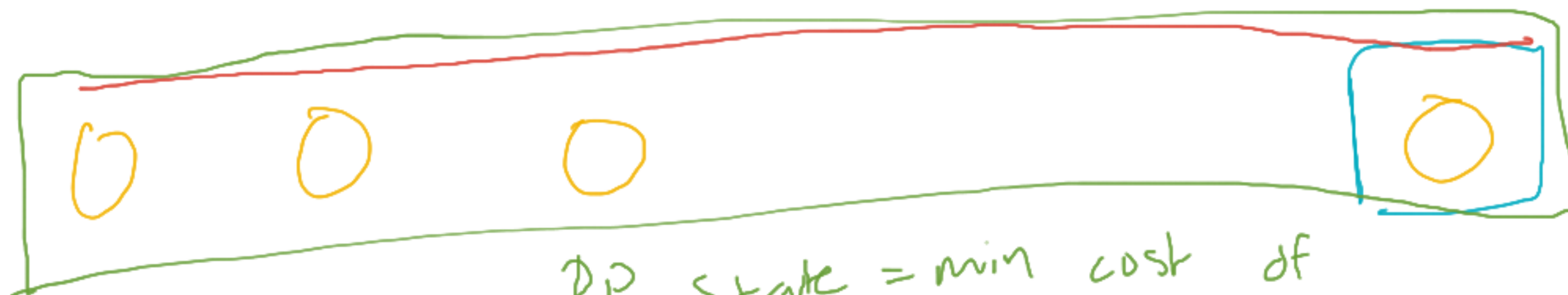
min cost of covering cows $i \dots k$

Larger umbrellas do not necessarily cost more than smaller umbrellas.



2 → 8
3 → 7

$$c'(l) = \min_{l' \geq l} c(l')$$



DP State = min cost of
covering all of the