

Discrete Optimization Problem

$$U = \{u_1, u_2, \dots, u_n\}$$

$$\forall S \subseteq U$$

$$f: S \rightarrow \mathbb{R}$$

Objective:

$$\max f(S)$$

s.t.

Constraints

e.g.

max - vertex cover

$$V = \{v_1, \dots, v_n\}$$

$$\forall S \subseteq V$$

$$\text{cov}(S) = \# \text{ edges } S \text{ hit}$$

$$\max \text{cov}(S)$$

s.t.

$$|S| \leq K$$

Monotonicity:

f is monotonic, if

$$f(S \cup \{x\}) \geq f(S)$$

e.g.

max-vertex-cover

$\text{cov}(S)$ is monotonic

Submodularity:

f is submodular, if adding an element to a subset of S has higher benefit than adding it to S^{\supseteq}

$$\forall S \subseteq U, \forall T \subseteq S: f(S \cup \{x\}) - f(S) \leq f(T \cup \{x\}) - f(T)$$

e.g. 1: U : A set of Balls, each w/ a Color

$f(S)$: The # Colors in S

$$S = \{ \text{orange}, \text{purple}, \text{brown}, \text{blue} \}$$

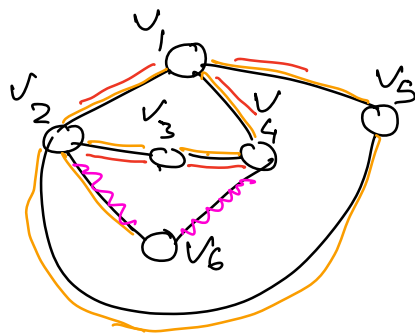
$$T = \{ \text{orange}, \text{blue}, \text{brown} \}$$

$$f(S \cup \{x\}) - f(S) \leq f(T \cup \{x\}) - f(T)$$

e.g. 2, max-Vertex-Cover:

is Cov Submodular?

Yes;



$$S = \{ v_1, v_2, v_3 \}$$

$$f(S) = 7$$

$$T = \{ v_1, v_3 \}$$

$$f(T) = 5$$

$$f(S \cup \{v_6\}) - f(S) = 1$$

$$f(T \cup \{v_6\}) - f(T) = 2$$

Theorem: if the objective function f is

① **monotonic**, and ② **Submodular**, then

Ⓐ the greedy approach satisfies a **$(1 - 1/e)$** approx. ratio
and ⑦ no other P-alg. can do better, unless $P=NP$.

* $(1 - 1/e)$ is a deviation of approx. ratio definition.

$\frac{1}{1 - 1/e}$ is the standard form

Max-Cover Problem

$$\text{Max } \left| \bigcup_{S_i \in \mathcal{S}} S_i \right| = f(x)$$

s.t.

$$|\mathcal{S}| \leq k$$

← Coverage

Set-Cover Problem

$$\text{Min } |\mathcal{S}|$$

s.t.

$$\left| \bigcup_{S_i \in \mathcal{S}} S_i \right| = |U|$$

Approx. ratio for Greedy
is $O(\log n)$

① $f(x)$ is monotonic: adding a new set to \mathcal{S} will not reduce $|\bigcup S_i|$

② $f(x)$ is Submodular:

$$\forall T \subseteq \mathcal{S} \subseteq U$$

$$f(T \cup \{x\}) - f(T) \geq f(\mathcal{S} \cup \{x\}) - f(\mathcal{S})$$

$$\left| \bigcup_{S_i \in T \cup \{x\}} S_i \right| - |T| \geq \left| \bigcup_{S_i \in \mathcal{S} \cup \{x\}} S_i \right| - |\mathcal{S}|$$

✓ Case 1: $x \in T \Rightarrow \left| \bigcup_{S_i \in T \cup \{x\}} S_i \right| = |T| \Rightarrow \left| \bigcup_{S_i \in T \cup \{x\}} S_i \right| - |T| = 0$

$\left| \bigcup_{S_i \in \mathcal{S} \cup \{x\}} S_i \right| = |\mathcal{S}| \Rightarrow \left| \bigcup_{S_i \in \mathcal{S} \cup \{x\}} S_i \right| - |\mathcal{S}| = 0$

$$\checkmark \text{Case 2: } x \notin T, x \in U \Rightarrow \begin{aligned} |\overset{T \cup \{x\}}{U}| - |T| &= 1 \\ |\overset{S \cup \{x\}}{U}| - |S| &= 0 \end{aligned}$$

$$\checkmark \text{Case 3: } x \notin T, x \notin U \Rightarrow \begin{aligned} |\overset{T \cup \{x\}}{U}| - |T| &= 1 \\ |\overset{S \cup \{x\}}{U}| - |S| &= 1 \end{aligned}$$

\Rightarrow Greedy Satisfies $(1 - 1/e)$ -approx. ratio for Max-Cover Problem.