

LP-Relaxation

IP
Min $f(x)$

$$Ax \leq b$$

$$\forall x_i \in X$$

$$x_i \in \mathbb{Z}$$

NP-Complete

$$f(x^+)$$

LP

Min $f(x)$
s.t.

$$Ax \leq b$$

$$\forall x_i \in X$$

$$x_i \in \mathbb{R}$$

P

$$f(x^*)$$

Similarly, for Max, LP provides an upper bound on the opt. solution of IP

LP Relaxation

Step1: Model your problem as IP

Step2: Relax IP to LP

Step3: $x^* \leftarrow \text{Solve (LP)}$

Step4: "Round" x^* to integer numbers x^+

VC-approximation Using LP-relaxation

V-C

v_1, \dots, v_n : vertices

e_1, \dots, e_m : edges

Let $x_i \in \{0, 1\}$ be 1 if v_i is selected, 0 otherwise

Min $\sum x_i$
S.t.
 $\forall (v_i, v_j) \in E$
 $x_i + x_j \geq 1$
 $\forall i \in [1, n]$
 $x_i \in \{0, 1\}$

Step 1 \rightarrow

Step 2 Relax to LP \rightarrow

Min $\sum x_i$
S.t.
 $\forall (v_i, v_j) \in E$
 $x_i + x_j \geq 1$
 $\forall i$
 $x_i \in [0, 1]$

Step 3
Solve (LP) \rightarrow

Step 4: Rounding x^* is the optimal values for LP

$x_i^+ = \text{round}(x_i^*)$ to closest integer

\nwarrow approx. solution

Correctness

$\forall (v_i, v_j)$ $x_i^* + x_j^* \geq 1 \Rightarrow$ at least one of x_i^* or x_j^* should be larger (or equal) to 0.5

\Rightarrow at least one of x_i^+ or x_j^+ should be 1

LP-relaxation provides a 2-approx. alg. for v-c.

$$\text{OPT} \geq \sum x_i^* \quad (\text{The opt. solution for LP})$$

$$= \sum_{x_i^* < 0.5} x_i^* + \sum_{x_j^* \geq 0.5} x_j^*$$

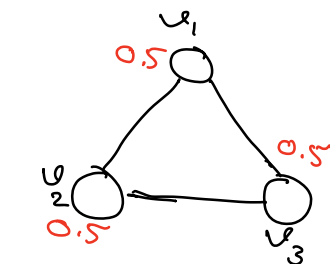
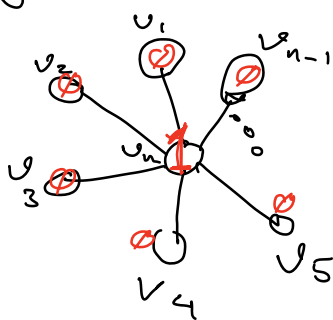
$$\geq \sum_{x_j^* \geq 0.5} x_j^*$$

$$\geq \sum_{x_j^* \geq 0.5} 0.5 = 0.5 \sum_{x_j^* \geq 0.5} 1 = 0.5 \|A\|$$

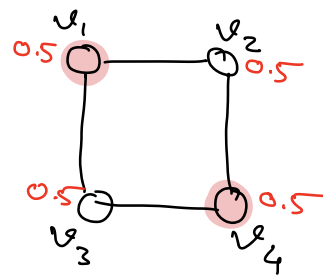
$$\Rightarrow \text{OPT} \geq 0.5 \|A\|$$

$$\Rightarrow \frac{\|A\|}{\text{opt}} \leq 2$$

eg.



approx: $\{v_1, v_2, v_3\}$
 opt: $\{v_1, v_2\}$
 $3/2$



approx: $\{v_1, v_2, v_3, v_4\}$
 opt: $\{v_1, v_4\}$

$4/2 = 2$
 → approx. ratio is Tight

Weighted Vertex Cover :

- every vertex is associated w/ a weight

- Goal: $\min \sum w_i x_i$

LP-relaxation:

Change the obj. function to
 $\min \sum w_i x_i$