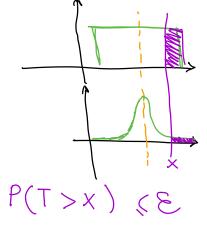


we want to design alg. that with a high Probe their Runtime is less than a given value Tail Probe



Occupancy Problems

No balls 1 2 000 1 bins

-given n balls, throw the balls in the bins independently.

⇒ Expected # balls in each bin = 1

Clain: The Prob. of having more than log(n) balls of in any bin is very low.

Proof:

P(having exactly k balls in bin i) = P(Bi=k)

 $P(s_i = k) = \binom{n}{k} \binom{n}{n}^k \left(\frac{n-1}{n}\right)^{n-k}$

 $\binom{n}{k}^{k} \leqslant \binom{n}{k} \leqslant \left(\frac{ne}{k}\right)^{k}$

 $\Rightarrow P(B_i = k) \leq \left(\frac{ne}{k}\right)^k {\binom{1}{n}}^k$ $= {\binom{e}{k}}^k$

Using Urion Bound

 $P(B_i \geq k) \leqslant \sum_{i=k}^{n} (e_{i})^{i}$

$$P(\tilde{y}_{\epsilon_{i}}) = P(\epsilon_{i}) + P(\epsilon_{2}) + \dots + P(\epsilon_{n})$$

$$- P(\epsilon_{i} \cap \epsilon_{2}) - P(\epsilon_{i} \cap \epsilon_{3}) - \dots - P(\epsilon_{n-1} \epsilon_{n})$$

$$+ P(\epsilon_{i} \cap \epsilon_{2} \cap \epsilon_{3}) - \dots$$

$$P(B_{i} \ge k) \le \sum_{i=k}^{n} (e_{i})^{i} \le \sum_{i=k}^{n} (e_{k})^{i}$$

$$= (e_{k})^{k} + (e_{k})^{k+1} + \dots + (e_{k})^{n}$$

$$\le (e_{k})^{k} (1 + e_{k})^{k} + (e_{k})^{2} + \dots)$$

$$\le (e_{k})^{k} (1 + e_{k})^{2} + \dots$$

$$\le (e_{k})^{k} (1 + e_{k})^{2} + \dots$$

K>@

Using Wolfram Alpha, if K=O(logn) = [(31nn)/Inlnn] $P(B_i > k) \leq {\binom{e_k}{k}}^k \frac{1}{1 - e_k}$

$$P(B_{i}^{*} \geq k) \leq ?$$

$$B_{i}^{*} = \max_{i \geq 1} R_{i}$$

$$P(B_{i}^{*} \geq k) = P(\bigcup B_{i} \geq k)$$

$$\leq \sum_{i \in 1}^{n} P(B_{i} \geq k)$$

=> P(having a bin with more than lug(n) balls) & 1/2