

Bin Packing (BP)

Given $U = \{w_1, \dots, w_n\}$, $0 < w_i \leq 1$,
what is the min # bins each
of max capacity $\frac{1}{2}$ to pack all
the elements. ✓

e.g.:

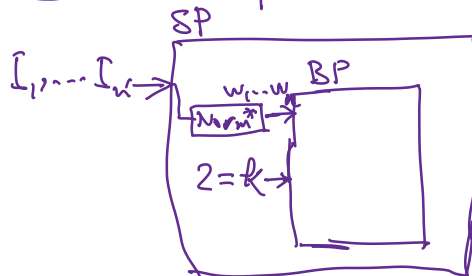
0.2	0.3	0.8	0.9	0.6
I_1	I_2			I_6

I_1, I_4
I_2, I_5
I_3, I_6

BP \in NP-Complete

① BP \in NP ✓

② $SP \leq_p BP$



$$w_j = \frac{2 I_j}{\sum_{l=1}^n I_l}$$

0/1 Knapsack

Given n Items, each with Profit P_i and weight w_i , and a
backpack with capacity C ,
what is the set S of items
to select s.t.

① $\sum_{I_j \in S} w_j \leq C$

② $\sum_{I_j \in S} P_j$ is Maximized

Dynamic Programming: $O(nC)$

$$Z = \log^C$$

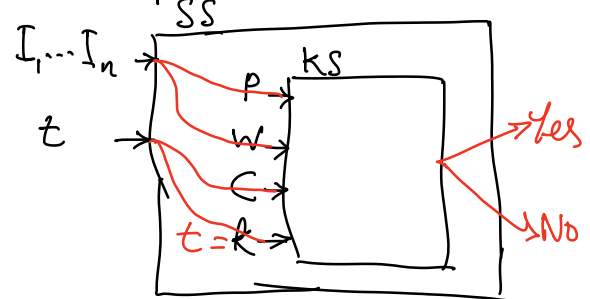
$$O(nC) = O(n2^Z)$$

exponential to the Input
Size
Pseudo Polynomial

0/1 Knapsack \in NP-Complete
(KS)

① KS \in NP ✓

② $SS \leq_p KS$



0/1 (Binary) Integer Programming

Given a Matrix $A_{n \times m}$
and a vector of binary
variables X , and a
vector of values b ,
find out if there exists
an assignment to X
Such that

$$AX \leq b$$

$$\begin{cases} x_1 + 3x_2 + 5x_3 \leq 2 \\ -2x_1 + 5x_2 + 4x_3 \leq 10 \\ x_1 - x_2 - x_3 \leq -5 \\ x_1, x_2, x_3 \in \{0, 1\} \end{cases}$$

$$A = \begin{bmatrix} 1 & 3 & 5 \\ -2 & 5 & 4 \\ 1 & -1 & -1 \end{bmatrix}$$

$$x = \langle x_1, x_2, x_3 \rangle$$

$$b = \langle 2, 10, -5 \rangle$$

$$Ax \leq b$$