DBMS Assignment 2

Team:

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Part 1

Consider a relation schema R(A, B, C, D, E, G) and the following sets of functional dependencies

```
F1 = {ABC \rightarrow D, BC \rightarrow EA, BCE \rightarrow G}
F2 = {A\rightarrowB, C\rightarrowAD, AE\rightarrowCG, BC \rightarrow C}
F3 = {AC \rightarrow B, BC \rightarrow D, BD \rightarrow E, AE \rightarrow G, ED \rightarrow A, DA \rightarrow C}
```

A. For each set F1, F2 and F3:-

Work out whether relation R with the respective set is in BCNF and show how you reached the answer.

Solution:

1) **F1**: True

Candidate key -> {BC}

1NF - True

2NF - True (The table should not have partial dependency on the candidate key)

3nf - True (No transitive dependency as BC is at the left side of all equations)

BCNF - True (All equations satisfy the criteria super key -> non prime attributes)

2) F2: False

Candidate Key -> {CE, AE}

1NF -> True

2NF -> In relation C -> AD, D is dependent on a subset of the candidate key(CE) hence not in 2nf

BCNF -> Not in BCNF as it is not in 2NF and criteria super key -> non prime attributes not satisfied for relation A -> B

3) **F3**: False

Candidate Key -> {AC},{ED},{DA},{BC},{BD}

1NF -> True

2NF -> True

3NF -> True

BCNF -> False, Not in BCNF (AE->G) violates the condition as AE is not a candidate key

B. For each set F1, F2 and F3:-

Find all candidate keys of R and show how you reached the answer.

```
Solution:
```

```
1) F1 ->
```

Given: F1 = {ABC \rightarrow D, BC \rightarrow EA, BCE \rightarrow G}

Attributes that can be determined (Are on the right side of the relation): ADEG

Attributes that cannot be determined: BC

Hence we start by checking BC

{BC} -> {BCEA}, now we check BCEA

{BCEA} -> {BCEADG}

Hence BC is a candidate key.

Hence Candidate Key: {BC}

2) F2 ->

Given: F2 = { $A \rightarrow B$, $C \rightarrow AD$, $AE \rightarrow CG$, $BC \rightarrow C$ }

Attributes that can be determined (Are on the right side of the relation): ABCDG

Attributes that cannot be determined: E

Hence we start by checking E

{E} -> Cannot determine anything alone

We now check AE

{AE} -> {AECGB}

{AECG} -> { AECGBD} Hence AE is a candidate key

Now we check CE

{CE} -> {CEAD}

{CEAD} -> {CEADBG} Hence CE is also a candidate key

Hence Candidate Key : {AE,CE}

3) F3 ->

Given: F3 = {AC
$$\rightarrow$$
 B, BC \rightarrow D, BD \rightarrow E, AE \rightarrow G, ED \rightarrow A, DA \rightarrow C}

Since all the attributes can be determined (are on the right hand side of the relations) we check the candidate keys manually.

{AC} -> {ABCDG}

{ED} -> {ABCDG}

{DA} -> {ABCDG}

{BC} -> {ABCDG}

{BD} -> {ABCDG}

Hence Candidate Key: {AC},{ED},{DA},{BC},{BD}

C. For each set F1, F2 and F3:-

Consider partitioning R into 3 sub relations $R1\{A,B,C\}$, $R2\{D,E,G\}$, $R3\{B,C,D\}$. Is this decomposition lossless? Explain your answer.

1) Given: F1 = {ABC \rightarrow D, BC \rightarrow EA, BCE \rightarrow G}

Check if R1 2 R3. R1 2 R3 = {BC}.

BC is a candidate key of both R1 and R3. Hence lossless decomposition

Check R13 and R2. R13 2 R2 = {D}

D is not the candidate key of R2 or R13. So the decomposition is lossy

R1 \square R2. R1 \square R2 = $\{\varnothing\}$. As it is a null set, the decomposition is lossy

Checking R2, R3. R2 2 R3 = {D}

As D is not a candidate key of either R3 or R2. We can say that the decomposition is lossy.

Since all individual decompositions are lossy, we can say that the decomposition R1,R2,R3 is lossy.

2) Given: F2 = {A \rightarrow B, C \rightarrow AD, AE \rightarrow CG, BC \rightarrow C} Candidate key = {AE,CE}

First, we check R1 $\$ 2 R2. R1 $\$ 2 R2 = $\{\emptyset\}$. As the intersection of R1 and R2 is a null set, the decomposition is lossy.

Now, we check R1 and R3. R1 \square R3 = {BC}.

 $\{BC\}+=\{A, B, C, D\}$. We can see that BC is candidate key of both R1 and R3. Hence the decomposition R1 and R3 is lossless.

Now, we check R13 = {A, B, C, D} and R2. R13 2 R2 = {D}

 $\{D\}+=\{D\}$. Thus, we can see that D is not a candidate key of either R13 or R2. So, the decomposition R13 and R2 is lossy.

 $\{D\}+=\{D\}$. Thus, we can see that D is not a candidate key of either R3 or R2. Hence the decomposition R2 and R3 is lossy.

Since all individual decompositions are lossy, the decomposition R1, R2 and R3 is lossy.

3) Given: $F3 = \{AC \rightarrow B, BC \rightarrow D, BD \rightarrow F, AF \rightarrow G, FD \rightarrow A, DA \rightarrow C\}$

3) Given: F3 = {AC \rightarrow B, BC \rightarrow D, BD \rightarrow E, AE \rightarrow G, ED \rightarrow A, DA \rightarrow C} Firstly, check R1 \bigcirc R3. R1 \bigcirc R3 = {BC}.

 $\{BC\}+=\{A, B, C, D, E, G\}$. BC is candidate key of both R1 and R3. So, the decomposition R1 and R3 is lossless.

Checking R1 $\$ 2 R2. R1 $\$ 2 R2 = $\$ 8. As the intersection of R1 and R2 is a null set, the decomposition is lossy.

To check R13 and R2

R13 = {A, B, C, D} and R2. R13 2 R2 = {D}

 $\{D\}+=\{D\}$. We can see that D is not a candidate key of either R13 or R2. So, the decomposition R13 and R2 is lossy.

 $\{D\}$ + = $\{D\}$. Thus, we can see that D is not a candidate key of either R3 or R2. Hence the decomposition R2 and R3 is lossy.

Since all individual decompositions are lossy, the decomposition R1, R2 and R3 is lossy.

Part 2

Entities:

1. Person(Person ID, Name, DOB, Gender)

Non-null attributes: Person_ID, Name, DOB, Gender

Key: Person_ID

IND: None

2. Employee(Person ID, Schedule, Employee Type, Salary per hour)

Non-null attributes: Person ID, Schedule, Employee Type, Salary per hour

Key: Person ID

IND: Employee[Person ID] ⊆ Person[Person ID]

3. Entry Log(Person ID, Timestamp)

Non-null attributes: Person ID, Timestamp

Key: Person ID, Timestamp

IND: Entry_Log[Person_ID] \subseteq Person[Person_ID]

4. Employee Exit log(Person ID, Timestamp)

Non-null attributes: Person ID, Timestamp

Key of Entry_Log: Person_ID, Timestamp

IND: Employee Exit log[Person ID] ⊆ Employee[Person ID]

5. Trainer(Person ID, Credentials)

Non null attributes: Person_ID, Credentials

Key: Person_ID

IND: Trainer[Person ID] ⊆ Employee[Person ID]

6. DeskEmployee(Person ID)

Non null attributes: Person ID

Key: Person ID

IND: DeskEmployee[Person ID] ⊆ Employee[Person ID]

7. Member(Person ID, Membership ID)

Non null attributes: Person ID, Membership_ID

Key: Person ID

IND: Member[Person ID] ⊆ Person[Person ID]

8. Family(Person ID, CreditCard)

Non null attributes: Person ID, CreditCard

Key: Person_ID

IND: Family[Person ID] ⊆ Member[Person ID]

9. University Affiliate(Person ID, Department)

Non null attributes: Person ID, Department

Key: Person_ID

IND: University_Affiliate[Person_ID] \subseteq Member[Person_ID]

10. Student(Person ID, Student Type)

Non null attributes: Person ID, Student Type

Key: Person_ID

IND: Student[Person ID] ⊆ University Affiliate[Person ID]

11. Non-Student(<u>Person_ID</u>, Member_type, Credit_card)

Non null attributes: Person_ID, Member_type, Credit_card

Key: Person ID

IND: Non-Student[Person_ID] \subseteq University_Affiliate[Person_ID]

12. Space(Space ID, Description, Max Capacity)

Non-null attributes: Space_ID, Description, Max_Capacity

Key: Space_ID

IND (INclusion Dependencies): None

13. Events(Event ID, Description, Start time, end time, capacity)

Non-null attributes: Event ID, Description, Start time, end time, capacity

Key: Event ID

IND (INclusion Dependencies): None

14. Location Reading(Person ID, Space ID, Sensor ID, Timestamp)

Key: Person ID, Space ID, Sensor ID, Timestamp

Non-null attributes: Person ID, Space ID, Sensor ID, Timestamp

 $\mathsf{IND}: \ \mathsf{Location_Reading}[\mathsf{Person_ID}] \subseteq \mathsf{Person}[\mathsf{Person_ID}]$

Location_Reading[Space_ID] \subseteq Space[Space_ID]

Location_Reading[Sensor_ID] \subseteq LocationSensor[Sensor_ID]

15. Equipment (Equipment ID, Equipment type, is available)

Key: Equipment ID

Non-null attributes: Equipment_ID, Equipment_type, is_available

IND (INclusion Dependencies): None

16. LocationSensor(Sensor ID, Coverage)

Non-null attributes: Sensor ID, Coverage

Key: Sensor_ID

IND (INclusion Dependencies): None

17. EquipmentSensor(Sensor_ID, Coverage)

Non-null attributes: Sensor_ID, Coverage

Key: Sensor_ID

IND (INclusion Dependencies): None

18. UsageReading(Equipment ID, Person ID, timestamp, Sensor ID)

Key of UsageReading: Equipment_ID, Person_ID, timestamp, Sensor_ID

Non-null attributes: Equipment_ID, Person_ID, timestamp, Sensor_ID

IND: UsageReading[Equipment_ID] ⊆ Equipment[Equipment_ID]

UsageReading[Person ID] ⊆ Person[Person ID]

UsageReading[Sensor_ID] ⊆ EquipmentSensor[Sensor_ID]

Relations:

1. MemberAttendsEvents(Person ID, Event ID)

Key: Person ID, Event ID

Non-null attributes: Person ID, Event ID

IND (INclusion Dependencies):

MemberAttendsEvents[Person ID] ⊆ Member[Person ID]

MemberAttendsEvents[Event ID] ⊆ Event[Event ID]

2. EventsHostedIn(EventId, Space Id)

Key: Event_id

Non-null attributes: Event id, Space Id

IND:

 ${\sf EventsHostedIn[Event_Id]} \subseteq {\sf Event[Event_Id]}$

Event[Event Id] ⊆ EventsHostedIn[Event Id]

EventsHostedIn[Space ID] ⊆ Space[Space ID]

3. SpaceContainsEquipments(Equipment Id, Space Id)

Key: Equipment Id

Non-null attributes: Equipment_Id, Space_Id

IND:

 $SpaceContainsEquipments[Equipment_Id] \ \subseteq \ Equipment[Equipment_Id]$

Equipment[Equipment Id] ⊆ SpaceContainsEquipments[Equipment Id]

SpaceContainsEquipments[Space ID] ⊆ Space[Space ID]

4. FamilyRealtedTo(<u>FamilyPersonId</u>, UniversityAffiliatePersonId)

Key: FamilyPersonId

Non-null attributes: FamilyPersonId, UniversityAffiliatePersonId

IND:

```
FamilyRealtedTo[FamilyPersonId] ⊆ Family[Person_ID]
Family[Person_ID] ⊆ FamilyRealtedTo[FamilyPersonId]
FamilyRealtedTo[UniversityAffiliatePersonId] ⊆ University_Affiliate[Person_ID]
```

Queries:

Entities

```
CREATE TABLE Person (
Person_ID INT PRIMARY KEY,
Name VARCHAR(255) NOT NULL,
DOB DATE NOT NULL,
Gender VARCHAR(10) NOT NULL
);

CREATE TABLE Employee (
Person_ID INT PRIMARY KEY,
Schedule VARCHAR(255) NOT NULL,
Employee_Type VARCHAR(255) NOT NULL,
Salary_per_hour DECIMAL(10,2) NOT NULL,
FOREIGN KEY (Person_ID) REFERENCES Person(Person_ID)
);
```

```
CREATE TABLE Entry Log (
  Person ID INT,
  Timestamp TIMESTAMP NOT NULL,
  FOREIGN KEY (Person ID) REFERENCES Person(Person ID)
);
CREATE TABLE Employee Exit Log (
  Person ID INT,
  Timestamp TIMESTAMP NOT NULL,
  FOREIGN KEY (Person_ID) REFERENCES Employee(Person_ID)
);
CREATE TABLE Trainer (
  Person ID INT PRIMARY KEY,
  Credentials VARCHAR(255) NOT NULL,
  FOREIGN KEY (Person_ID) REFERENCES Employee(Person_ID)
);
CREATE TABLE DeskEmployee (
  Person ID INT PRIMARY KEY,
  FOREIGN KEY (Person_ID) REFERENCES Employee(Person_ID)
);
CREATE TABLE Member (
  Person ID INT PRIMARY KEY,
  Membership ID INT NOT NULL,
  FOREIGN KEY (Person ID) REFERENCES Person(Person ID)
);
CREATE TABLE Family (
  Person ID INT PRIMARY KEY,
  CreditCard VARCHAR(255) NOT NULL,
  FOREIGN KEY (Person ID) REFERENCES Member(Person ID)
);
CREATE TABLE University Affiliate (
  Person ID INT PRIMARY KEY,
  Department VARCHAR(255) NOT NULL,
  FOREIGN KEY (Person_ID) REFERENCES Member(Person_ID)
);
CREATE TABLE Student (
  Person ID INT PRIMARY KEY,
  Student Type VARCHAR(255) NOT NULL,
```

```
FOREIGN KEY (Person ID) REFERENCES University Affiliate(Person ID)
);
CREATE TABLE Non Student (
  Person ID INT PRIMARY KEY,
  Member type VARCHAR(255) NOT NULL,
  Credit card VARCHAR(255) NOT NULL,
  FOREIGN KEY (Person ID) REFERENCES University Affiliate(Person ID)
);
CREATE TABLE Space (
  Space ID INT PRIMARY KEY,
  Description VARCHAR(255) NOT NULL,
  Max Capacity INT NOT NULL
);
CREATE TABLE Events (
  Event ID INT PRIMARY KEY,
  Description VARCHAR(255) NOT NULL,
  Start time DATETIME NOT NULL,
  End time DATETIME NOT NULL,
  Capacity INT NOT NULL
);
CREATE TABLE LocationSensor (
  Sensor ID INT PRIMARY KEY,
  Coverage VARCHAR(255) NOT NULL
);
CREATE TABLE Location Reading (
  Person ID INT,
  Space ID INT,
  Sensor ID INT,
  Timestamp TIMESTAMP NOT NULL,
  FOREIGN KEY (Person ID) REFERENCES Person(Person ID),
  FOREIGN KEY (Space ID) REFERENCES Space(Space ID),
  FOREIGN KEY (Sensor ID) REFERENCES LocationSensor(Sensor ID)
);
CREATE TABLE Equipment (
  Equipment ID INT PRIMARY KEY,
  Equipment type VARCHAR(255) NOT NULL,
  is available BOOLEAN NOT NULL
);
```

```
CREATE TABLE EquipmentSensor (
  Sensor ID INT PRIMARY KEY,
  Coverage VARCHAR(255) NOT NULL
);
CREATE TABLE UsageReading (
  Equipment ID INT,
  Person ID INT,
  Timestamp TIMESTAMP NOT NULL,
  Sensor ID INT,
  FOREIGN KEY (Equipment ID) REFERENCES Equipment(Equipment ID),
  FOREIGN KEY (Person ID) REFERENCES Person(Person ID),
  FOREIGN KEY (Sensor ID) REFERENCES EquipmentSensor(Sensor ID)
);
Relations
CREATE TABLE MemberAttendsEvents (
  Person ID INT,
  Event ID INT,
  PRIMARY KEY (Person ID, Event ID),
  FOREIGN KEY (Person ID) REFERENCES Member(Person ID),
  FOREIGN KEY (Event ID) REFERENCES Events(Event ID)
);
CREATE TABLE EventsHostedIn (
  Event ID INT PRIMARY KEY,
  Space ID INT,
  FOREIGN KEY (Event ID) REFERENCES Events(Event ID),
  FOREIGN KEY (Space ID) REFERENCES Space(Space ID)
);
CREATE TABLE SpaceContainsEquipments (
  Equipment ID INT PRIMARY KEY,
  Space ID INT,
  FOREIGN KEY (Equipment ID) REFERENCES Equipment(Equipment ID),
  FOREIGN KEY (Space_ID) REFERENCES Space(Space_ID)
);
CREATE TABLE FamilyRealtedTo (
```

```
FamilyPersonId INT PRIMARY KEY,
UniversityAffiliatePersonId INT,
FOREIGN KEY (FamilyPersonId) REFERENCES Family(Person_ID),
FOREIGN KEY (UniversityAffiliatePersonId) REFERENCES
University_Affiliate(Person_ID)
);
```