

On the Formalization of Martingales

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theory <i>Measure-Space-Addendum</i>		
imports <i>HOL-Analysis.Measure-Space</i>		
begin		

1 Sigma Algebra Generated by a Family of Functions

definition *sigma-gen* :: 'a set \Rightarrow 'b measure \Rightarrow ('a \Rightarrow 'b) set \Rightarrow 'a measure **where**
sigma-gen Ω N $S \equiv \text{sigma } \Omega (\bigcup f \in S. \{f - 'A \cap \Omega \mid A. A \in N\})$

lemma

shows *sets-sigma-gen*: *sets (sigma-gen Ω N S) = sigma-sets $\Omega (\bigcup f \in S. \{f - 'A \cap \Omega \mid A. A \in N\})$*

and *space-sigma-gen[simp]*: *space (sigma-gen Ω N S) = Ω*

by (*auto simp add: sigma-gen-def sets-measure-of-conv space-measure-of-conv*)

lemma *measurable-sigma-gen*:

assumes $f \in S$ $f \in \Omega \rightarrow \text{space } N$

shows $f \in \text{sigma-gen } \Omega$ N $S \rightarrow_M N$

using *assms* **by** (*intro measurableI, auto simp add: sets-sigma-gen*)

lemma *measurable-sigma-gen-singleton*:

assumes $f \in \Omega \rightarrow \text{space } N$

shows $f \in \text{sigma-gen } \Omega$ N $\{f\} \rightarrow_M N$

using *assms measurable-sigma-gen* **by** *blast*

lemma *measurable-iff-contains-sigma-gen*:

shows $(f \in M \rightarrow_M N) \iff f \in \text{space } M \rightarrow \text{space } N \wedge \text{sigma-gen } (\text{space } M) N$
 $\{f\} \subseteq M$

proof (*standard, goal-cases*)

case 1

hence $f \in \text{space } M \rightarrow \text{space } N$ **using** *measurable-space* **by** *fast*

thus ?case **unfolding** *sets-sigma-gen* **by** (*simp, intro sigma-algebra.sigma-sets-subset,*
(blast intro: sets.sigma-algebra-axioms measurable-sets[OF 1])+))

next

case 2

thus ?case **using** *measurable-mono[OF - refl - space-sigma-gen, of N M]* *measurable-sigma-gen-singleton* **by** *fast*

qed

```

lemma measurable-family-iff-contains-sigma-gen:
  shows  $(S \subseteq M \rightarrow_M N) \longleftrightarrow S \subseteq \text{space } M \rightarrow \text{space } N \wedge \text{sigma-gen } (\text{space } M)$ 
 $N \ S \subseteq M$ 
proof (standard, goal-cases)
  case 1
    hence subset:  $S \subseteq \text{space } M \rightarrow \text{space } N$  using measurable-space by fast
    have  $\{f \mid \neg A \cap \text{space } M \mid A. A \in N\} \subseteq M$  if  $f \in S$  for  $f$  using measurable-iff-contains-sigma-gen[unfolded sets-sigma-gen, of f] 1 subset that by blast
    then show ?case unfolding sets-sigma-gen using sets.sigma-algebra-axioms by
    (simp add: subset, intro sigma-algebra.sigma-sets-subset, blast+)
  next
    case 2
    hence subset:  $S \subseteq \text{space } M \rightarrow \text{space } N$  by simp
    show ?case
    proof (standard, goal-cases)
      case (1 x)
        have sigma-gen (space M) N  $\{x\} \subseteq M$  by (metis (no-types, lifting) 1 2
        sets-sigma-gen SUP-le-iff sigma-sets-le-sets-iff singletonD)
        thus ?case using measurable-iff-contains-sigma-gen subset[THEN subsetD, OF
        1] by fast
      qed
    qed

end
theory Elementary-Metric-Spaces-Addendum
imports HOL-Analysis.Elementary-Metric-Spaces
begin

```

2 Diameter Lemma

```

lemma diameter-comp-strict-mono:
  fixes  $s :: \text{nat} \Rightarrow 'a :: \text{metric-space}$ 
  assumes strict-mono  $r$  bounded  $\{s \ i \mid i. r \ n \leq i\}$ 
  shows diameter  $\{s \ (r \ i) \mid i. n \leq i\} \leq \text{diameter } \{s \ i \mid i. r \ n \leq i\}$ 
proof (rule diameter-subset)
  show  $\{s \ (r \ i) \mid i. n \leq i\} \subseteq \{s \ i \mid i. r \ n \leq i\}$  using assms(1) monotoneD
  strict-mono-mono by fastforce
qed (intro assms(2))

lemma diameter-bounded-bound':
  fixes  $S :: 'a :: \text{metric-space}$  set
  assumes  $S$ : bdd-above (case-prod dist '  $(S \times S)$ )  $x \in S \ y \in S$ 
  shows dist  $x \ y \leq \text{diameter } S$ 
proof -
  have  $(x, y) \in S \times S$  using  $S$  by auto
  then have dist  $x \ y \leq (\text{SUP } (x, y) \in S \times S. \text{dist } x \ y)$  by (rule cSUP-upper2[OF
  assms(1)]) simp
  with  $\langle x \in S \rangle$  show ?thesis by (auto simp: diameter-def)
qed

```

lemma *bounded-imp-dist-bounded*:
assumes *bounded* (*range s*)
shows *bounded* $((\lambda(i, j). \text{dist } (s \ i) \ (s \ j)) \ '(\{n..\} \times \{n..\}))$
using *bounded-dist-comp*[*OF* *bounded-fst* *bounded-snd*, *OF* *bounded-Times*(1,1)[*OF* *assms*(1,1)]] **by** (*rule* *bounded-subset*, *force*)

lemma *cauchy-iff-diameter-tends-to-zero-and-bounded*:
fixes $s :: \text{nat} \Rightarrow 'a :: \text{metric-space}$
shows $\text{Cauchy } s \longleftrightarrow ((\lambda n. \text{diameter } \{s \ i \mid i. i \geq n\}) \longrightarrow 0 \wedge \text{bounded } (\text{range } s))$
proof –
have $\{s \ i \mid i. N \leq i\} \neq \{\}$ **for** N **by** *blast*
hence *diameter-SUP*: $\text{diameter } \{s \ i \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s \ i) \ (s \ j))$ **for** N **unfolding** *diameter-def* **by** (*auto* *intro!*: *arg-cong*[*of* - - *Sup*])
show *?thesis*
proof $((\text{intro } \text{iffI}) ; \text{clarsimp})$
assume *asm*: *Cauchy s*
have $\exists N. \forall n \geq N. \text{norm } (\text{diameter } \{s \ i \mid i. n \leq i\}) < e$ **if** *e-pos*: $e > 0$ **for** e
proof –
obtain N **where** *dist-less*: $\text{dist } (s \ n) \ (s \ m) < (1/2) * e$ **if** $n \geq N \ m \geq N$
for $n \ m$ **using** *asm* *e-pos* **by** (*metis* *Cauchy-def* *mult-pos-pos* *zero-less-divide-iff* *zero-less-numeral* *zero-less-one*)
{
fix r **assume** $r \geq N$
hence $\text{dist } (s \ n) \ (s \ m) < (1/2) * e$ **if** $n \geq r \ m \geq r$ **for** $n \ m$ **using** *dist-less*
that **by** *simp*
hence $(\text{SUP } (i, j) \in \{r..\} \times \{r..\}. \text{dist } (s \ i) \ (s \ j)) \leq (1/2) * e$ **by** (*intro* *cSup-least*) *fastforce* +
also **have** $\dots < e$ **using** *e-pos* **by** *simp*
finally **have** $\text{diameter } \{s \ i \mid i. r \leq i\} < e$ **using** *diameter-SUP* **by** *presburger*
}
moreover **have** $\text{diameter } \{s \ i \mid i. r \leq i\} \geq 0$ **for** r **unfolding** *diameter-SUP*
using *bounded-imp-dist-bounded*[*OF* *cauchy-imp-bounded*, *THEN* *bounded-imp-bdd-above*, *OF* *asm*] **by** (*intro* *cSup-upper2*, *auto*)
ultimately **show** *?thesis* **by** *auto*
qed
thus $(\lambda n. \text{diameter } \{s \ i \mid i. n \leq i\}) \longrightarrow 0 \wedge \text{bounded } (\text{range } s)$ **using** *cauchy-imp-bounded*[*OF* *asm*] **by** (*simp* *add*: *LIMSEQ-iff*)
next
assume *asm*: $(\lambda n. \text{diameter } \{s \ i \mid i. n \leq i\}) \longrightarrow 0 \wedge \text{bounded } (\text{range } s)$
have $\exists N. \forall n \geq N. \forall m \geq N. \text{dist } (s \ n) \ (s \ m) < e$ **if** *e-pos*: $e > 0$ **for** e
proof –
obtain N **where** *diam-less*: $\text{diameter } \{s \ i \mid i. r \leq i\} < e$ **if** $r \geq N$ **for** r
using *LIMSEQ-D* *asm*(1) *e-pos* **by** *fastforce*
{
fix $n \ m$ **assume** $n \geq N \ m \geq N$
hence $\text{dist } (s \ n) \ (s \ m) < e$ **using** *cSUP-lessD*[*OF* *bounded-imp-dist-bounded*[*THEN* *bounded-imp-bdd-above*], *OF* *asm*(2)] *diam-less*[*unfolded* *diameter-SUP*]] **by** *auto*

```

    }
    thus ?thesis by blast
  qed
  then show Cauchy s by (simp add: Cauchy-def)
  qed
qed

end
theory Bochner-Integration-Addendum
  imports HOL-Analysis.Bochner-Integration Elementary-Metric-Spaces-Addendum
begin

```

3 Auxiliary Lemmas for Bochner Integration

3.1 Simple Functions

```

lemma integrable-implies-simple-function-sequence:
  fixes f :: 'a  $\Rightarrow$  'b :: {banach, second-countable-topology}
  assumes integrable M f
  obtains s where  $\bigwedge i. \text{simple-function } M (s\ i)$ 
    and  $\bigwedge i. \text{emeasure } M \{y \in \text{space } M. s\ i\ y \neq 0\} \neq \infty$ 
    and  $\bigwedge x. x \in \text{space } M \implies (\lambda i. s\ i\ x) \longrightarrow f\ x$ 
    and  $\bigwedge x\ i. x \in \text{space } M \implies \text{norm } (s\ i\ x) \leq 2 * \text{norm } (f\ x)$ 
proof -
  have f:  $f \in \text{borel-measurable } M (\int^+ x. \text{norm } (f\ x) \partial M) < \infty$  using assms
  unfolding integrable-iff-bounded by auto
  obtain s where s:  $\bigwedge i. \text{simple-function } M (s\ i) \bigwedge x. x \in \text{space } M \implies (\lambda i. s\ i\ x) \longrightarrow f\ x$ 
     $\bigwedge i x. x \in \text{space } M \implies \text{norm } (s\ i\ x) \leq 2 * \text{norm } (f\ x)$  using
    borel-measurable-implies-sequence-metric[OF f(1)] unfolding norm-conv-dist by
    metis
  {
    fix i
    have  $(\int^+ x. \text{norm } (s\ i\ x) \partial M) \leq (\int^+ x. \text{ennreal } (2 * \text{norm } (f\ x)) \partial M)$  using
    s by (intro nn-integral-mono, auto)
    also have  $\dots < \infty$  using f by (simp add: nn-integral-cmult ennreal-mult-less-top
    ennreal-mult)
    finally have sbi: Bochner-Integration.simple-bochner-integrable M (s i) using
    s by (intro simple-bochner-integrableI-bounded) auto
    hence  $\text{emeasure } M \{y \in \text{space } M. s\ i\ y \neq 0\} \neq \infty$  by (auto intro: integrableI-simple-bochner-integrable
    simple-bochner-integrable.cases)
  }
  thus ?thesis using that s by blast
qed

lemma simple-function-indicator-representation:
  fixes f :: 'a  $\Rightarrow$  'b :: {second-countable-topology, banach}
  assumes f: simple-function M f and x:  $x \in \text{space } M$ 
  shows  $f\ x = (\sum y \in f\ ^\text{'space } M. \text{indicator } (f\ -^\text{'}\{y\} \cap \text{space } M) x *_R y)$ 

```

(is ?l = ?r)
proof –
 have ?r = ($\sum y \in f \text{ ' space } M$.
 (if $y = f x$ then indicator ($f - \{y\} \cap \text{space } M$) $x *_R y$ else 0)) **by** (auto intro!:
 sum.cong)
 also have ... = indicator ($f - \{f x\} \cap \text{space } M$) $x *_R f x$ **using** assms **by** (auto
 dest: simple-functionD)
 also have ... = $f x$ **using** x **by** (auto simp: indicator-def)
 finally show ?thesis **by** auto
qed

lemma simple-function-indicator-representation-AE:
 fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
 assumes f : simple-function $M f$
 shows AE x in M . $f x = (\sum y \in f \text{ ' space } M$. indicator ($f - \{y\} \cap \text{space } M$) x
 $*_R y$)
by (metis (mono-tags, lifting) AE-I2 simple-function-indicator-representation f)

lemmas simple-function-scaleR[intro] = simple-function-compose2[**where** $h = (*_R)$]
lemmas integrable-simple-function = simple-bochner-integrable.intros[**THEN** has-bochner-integral-simple-boch
THEN integrable.intros]

lemma integrable-simple-function-induct[consumes 2, case-names cong indicator
 add, induct set: simple-function]:
 fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
 assumes f : simple-function $M f$ emeasure $M \{y \in \text{space } M. f y \neq 0\} \neq \infty$
 assumes cong: $\bigwedge g$. simple-function $M f \implies \text{emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty$
 $\implies \text{simple-function } M g \implies \text{emeasure } M \{y \in \text{space } M. g y \neq 0\} \neq \infty$
 $\implies (\bigwedge x. x \in \text{space } M \implies f x = g x) \implies P f \implies P g$
 assumes indicator: $\bigwedge A y. A \in \text{sets } M \implies \text{emeasure } M A < \infty \implies P (\lambda x.$
 indicator $A x *_R y)$
 assumes add: $\bigwedge f g$. simple-function $M f \implies \text{emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty \implies$
 simple-function $M g \implies \text{emeasure } M \{y \in \text{space } M. g y \neq 0\} \neq \infty \implies$
 $(\bigwedge z. z \in \text{space } M \implies \text{norm } (f z + g z) = \text{norm } (f z) + \text{norm } (g z)) \implies$
 $P f \implies P g \implies P (\lambda x. f x + g x)$
 shows $P f$
proof –
 let $?f = \lambda x. (\sum y \in f \text{ ' space } M$. indicat-real ($f - \{y\} \cap \text{space } M$) $x *_R y$)
 have f -ae-eq: $f x = ?f x$ **if** $x \in \text{space } M$ **for** x **using** simple-function-indicator-representation[OF
 $f(1)$ that] .
 moreover have emeasure $M \{y \in \text{space } M. ?f y \neq 0\} \neq \infty$ **by** (metis (no-types,
 lifting) Collect-cong calculation $f(2)$)

moreover have $P (\lambda x. \sum_{y \in S}. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y)$
 $\text{simple-function } M (\lambda x. \sum_{y \in S}. \text{indicat-real } (f - \{y\} \cap \text{space } M) x$
 $*_R y)$
 $\text{emeasure } M \{y \in \text{space } M. (\sum_{x \in S}. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \neq \infty$
if $S \subseteq f - \{ \text{space } M \}$ **for** S **using** $\text{simple-functionD}(1)[\text{OF } \text{assms}(1),$
 $\text{THEN } \text{rev-finite-subset, OF that}] \text{ that}$
proof (*induction rule: finite-induct*)
case empty
 $\{$
case 1
then show $?case \text{ using indicator}[of \{ \}] \text{ by force}$
next
case 2
then show $?case \text{ by force}$
next
case 3
then show $?case \text{ by force}$
 $\}$
next
case (*insert x F*)
have $(f - \{x\} \cap \text{space } M) \subseteq \{y \in \text{space } M. f y \neq 0\}$ **if** $x \neq 0$ **using that by**
 blast
moreover have $\{y \in \text{space } M. f y \neq 0\} = \text{space } M - (f - \{0\} \cap \text{space } M)$
by blast
moreover have $\text{space } M - (f - \{0\} \cap \text{space } M) \in \text{sets } M$ **using** $\text{simple-functionD}(2)[\text{OF } f(1)] \text{ by blast}$
ultimately have $\text{fin-0: emeasure } M (f - \{x\} \cap \text{space } M) < \infty$ **if** $x \neq 0$
using that by (*metis emeasure-mono f(2) infinity-ennreal-def top.not-eq-extremum top-unique*)
hence $\text{fin-1: emeasure } M \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R$
 $x \neq 0\} \neq \infty$ **if** $x \neq 0$ **by** (*metis (mono-tags, lifting) emeasure-mono f(1) indica-*
 $\text{tor-simps}(2) \text{ linorder-not-less mem-Collect-eq scaleR-eq-0-iff simple-functionD}(2)$
 $\text{subsetI that})$

have $*$: $(\sum_{y \in \text{insert } x F}. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y) = (\sum_{y \in F}. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y) + \text{indicat-real } (f - \{x\} \cap \text{space } M)$
 $x *_R x$ **for** x **by** (*metis (no-types, lifting) Diff-empty Diff-insert0 add commute insert.hyps(1) insert.hyps(2) sum.insert-remove*)
have $**$: $\{y \in \text{space } M. (\sum_{x \in \text{insert } x F}. \text{indicat-real } (f - \{x\} \cap \text{space } M) y$
 $*_R x) \neq 0\} \subseteq \{y \in \text{space } M. (\sum_{x \in F}. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x)$
 $\neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\}$ **unfolding**
 $*$ **by** fastforce
 $\{$
case 1
hence $x: x \in f - \{ \text{space } M \}$ **and** $F: F \subseteq f - \{ \text{space } M \}$ **by auto**
show $?case$
proof (*cases x = 0*)
case True

```

    then show ?thesis unfolding * using insert(3)[OF F] by simp
next
  case False
  have norm-argument: norm (( $\sum_{y \in F} \text{indicat-real } (f - \{y\} \cap \text{space } M) z$ 
 $*_R y) + \text{indicat-real } (f - \{x\} \cap \text{space } M) z *_R x) = \text{norm } (\sum_{y \in F} \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) + \text{norm } (\text{indicat-real } (f - \{x\} \cap \text{space } M) z *_R x)$ 
  if  $z: z \in \text{space } M$  for  $z$ 
  proof (cases  $f z = x$ )
  case True
  have indicat-real  $(f - \{y\} \cap \text{space } M) z *_R y = 0$  if  $y \in F$  for  $y$  using
  True insert(2)  $z$  that 1 unfolding indicator-def by force
  hence  $(\sum_{y \in F} \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) = 0$  by (meson
  sum.neutral)
  then show ?thesis by force
  next
  case False
  then show ?thesis by force
qed
show ?thesis using False simple-functionD(2)[OF  $f(1)$ ] insert(3,5)[OF F]
simple-function-scaleR fin-0 fin-1 by (subst *, subst add, subst simple-function-sum)
(blast intro: norm-argument indicator)+
qed
next
  case 2
  hence  $x: x \in f \text{ ' space } M$  and  $F: F \subseteq f \text{ ' space } M$  by auto
  show ?case
  proof (cases  $x = 0$ )
  case True
  then show ?thesis unfolding * using insert(4)[OF F] by simp
  next
  case False
  then show ?thesis unfolding * using insert(4)[OF F] simple-functionD(2)[OF
 $f(1)$ ] by fast
  qed
next
  case 3
  hence  $x: x \in f \text{ ' space } M$  and  $F: F \subseteq f \text{ ' space } M$  by auto
  show ?case
  proof (cases  $x = 0$ )
  case True
  then show ?thesis unfolding * using insert(5)[OF F] by simp
  next
  case False
  have emeasure  $M \{y \in \text{space } M. (\sum_{x \in \text{insert } x F} \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \leq \text{emeasure } M \{y \in \text{space } M. (\sum_{x \in F} \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\}$ 
  using ** simple-functionD(2)[OF insert(4)[OF F]] simple-functionD(2)[OF
 $f(1)$ ] by (intro emeasure-mono, force+)

```


also have ... \leq *emeasure* $M \{y \in \text{space } M. (\sum x \in F. \text{indicat-real } (f - ' \{x\} \cap \text{space } M) y *_R x) \neq 0\} + \text{emeasure } M \{y \in \text{space } M. \text{indicat-real } (f - ' \{x\} \cap \text{space } M) y *_R x \neq 0\}$
using *simple-functionD*(2)[*OF insert*(4)[*OF F*]] *simple-functionD*(2)[*OF f*(1)] **by** (*intro emeasure-subadditive, force+*)
also have ... $< \infty$ **using** *insert*(5)[*OF F*] *fin-1*[*OF False*] **by** (*simp add: less-top*)
finally show ?thesis **by** *simp*
qed
}
qed
moreover have *simple-function* $M (\lambda x. \sum y \in f - ' \text{space } M. \text{indicat-real } (f - ' \{y\} \cap \text{space } M) x *_R y)$ **using** *calculation by blast*
moreover have $P (\lambda x. \sum y \in f - ' \text{space } M. \text{indicat-real } (f - ' \{y\} \cap \text{space } M) x *_R y)$ **using** *calculation by blast*
ultimately show ?thesis **by** (*intro cong[OF - - f(1,2)], blast, presburger+*)
qed

lemma *integrable-simple-function-induct-nn*[*consumes 3, case-names cong indicator add, induct set: simple-function*]:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes f : *simple-function* $M f$ *emeasure* $M \{y \in \text{space } M. f y \neq 0\} \neq \infty \wedge x. x \in \text{space } M \longrightarrow f x \geq 0$

assumes *cong*: $\bigwedge f g. \text{simple-function } M f \Longrightarrow \text{emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty \Longrightarrow (\bigwedge x. x \in \text{space } M \Longrightarrow f x \geq 0) \Longrightarrow \text{simple-function } M g \Longrightarrow \text{emeasure } M \{y \in \text{space } M. g y \neq 0\} \neq \infty \Longrightarrow (\bigwedge x. x \in \text{space } M \Longrightarrow g x \geq 0) \Longrightarrow (\bigwedge x. x \in \text{space } M \Longrightarrow f x = g x) \Longrightarrow P f \Longrightarrow P g$

assumes *indicator*: $\bigwedge A y. y \geq 0 \Longrightarrow A \in \text{sets } M \Longrightarrow \text{emeasure } M A < \infty \Longrightarrow P (\lambda x. \text{indicator } A x *_R y)$

assumes *add*: $\bigwedge f g. (\bigwedge x. x \in \text{space } M \Longrightarrow f x \geq 0) \Longrightarrow \text{simple-function } M f \Longrightarrow \text{emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty \Longrightarrow$

$(\bigwedge x. x \in \text{space } M \Longrightarrow g x \geq 0) \Longrightarrow \text{simple-function } M g \Longrightarrow \text{emeasure } M \{y \in \text{space } M. g y \neq 0\} \neq \infty \Longrightarrow$

$(\bigwedge z. z \in \text{space } M \Longrightarrow \text{norm } (f z + g z) = \text{norm } (f z) + \text{norm } (g z)) \Longrightarrow$

$P f \Longrightarrow P g \Longrightarrow P (\lambda x. f x + g x)$

shows $P f$

proof –

let ?f = $\lambda x. (\sum y \in f - ' \text{space } M. \text{indicat-real } (f - ' \{y\} \cap \text{space } M) x *_R y)$

have *f-ae-eq*: $f x = ?f x$ **if** $x \in \text{space } M$ **for** x **using** *simple-function-indicator-representation*[*OF f*(1) *that*] .

moreover have *emeasure* $M \{y \in \text{space } M. ?f y \neq 0\} \neq \infty$ **by** (*metis (no-types, lifting) Collect-cong calculation f(2)*)

moreover have $P (\lambda x. \sum y \in S. \text{indicat-real } (f - ' \{y\} \cap \text{space } M) x *_R y)$

simple-function $M (\lambda x. \sum y \in S. \text{indicat-real } (f - ' \{y\} \cap \text{space } M) x$

$*_R y)$

emeasure $M \{y \in \text{space } M. (\sum x \in S. \text{indicat-real } (f - ' \{x\} \cap \text{space$

$M) y *_R x) \neq 0\} \neq \infty$
 $\bigwedge x. x \in \text{space } M \implies 0 \leq (\sum_{y \in S. \text{indicat-real } (f - \{y\} \cap \text{space } M)$
 $x *_R y)$
if $S \subseteq f - \{y\} \cap \text{space } M$ **for** S **using** $\text{simple-functionD}(1)[\text{OF } \text{assms}(1),$
THEN $\text{rev-finite-subset, OF that}$ **that**
proof ($\text{induction rule: finite-subset-induct}$)
case empty
{
case 1
then show $?case$ **using** $\text{indicator}[of\ 0\ \{\}]$ **by force**
next
case 2
then show $?case$ **by force**
next
case 3
then show $?case$ **by force**
next
case 4
then show $?case$ **by force**
}
next
case ($\text{insert } x\ F$)
have $(f - \{x\} \cap \text{space } M) \subseteq \{y \in \text{space } M. f\ y \neq 0\}$ **if** $x \neq 0$ **using that by**
 blast
moreover have $\{y \in \text{space } M. f\ y \neq 0\} = \text{space } M - (f - \{0\} \cap \text{space } M)$
by blast
moreover have $\text{space } M - (f - \{0\} \cap \text{space } M) \in \text{sets } M$ **using** $\text{simple-functionD}(2)[\text{OF } f(1)]$ **by blast**
ultimately have $\text{fin-0: } \text{emeasure } M\ (f - \{x\} \cap \text{space } M) < \infty$ **if** $x \neq 0$
using that by ($\text{metis } \text{emeasure-mono } f(2)\ \text{infinity-ennreal-def top.not-eq-extremum}$
 top-unique)
hence $\text{fin-1: } \text{emeasure } M\ \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M)\ y *_R$
 $x \neq 0\} \neq \infty$ **if** $x \neq 0$ **by** ($\text{metis } (\text{mono-tags, lifting})\ \text{emeasure-mono } f(1)\ \text{indica-}$
 $\text{tor-simps}(2)\ \text{linorder-not-less mem-Collect-eq scaleR-eq-0-iff simple-functionD}(2)$
 subsetI that)

have $\text{nonneg-}x: x \geq 0$ **using** $\text{insert } f$ **by blast**
have $*$: $(\sum_{y \in \text{insert } x\ F. \text{indicat-real } (f - \{y\} \cap \text{space } M)\ x *_R y) =$
 $(\sum_{y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M)\ x *_R y) + \text{indicat-real } (f - \{x\} \cap$
 $\text{space } M)\ x *_R x$ **for** x **by** ($\text{metis } (\text{no-types, lifting})\ \text{add.commute insert.hyps}(1)$
 $\text{insert.hyps}(4)\ \text{sum.insert}$)
have $**$: $\{y \in \text{space } M. (\sum_{x \in \text{insert } x\ F. \text{indicat-real } (f - \{x\} \cap \text{space } M)\ y$
 $*_R x) \neq 0\} \subseteq \{y \in \text{space } M. (\sum_{x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M)\ y *_R x)$
 $\neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M)\ y *_R x \neq 0\}$ **unfolding**
 $*$ **by fastforce**
{
case 1
show $?case$
proof ($\text{cases } x = 0$)

```

    case True
    then show ?thesis unfolding * using insert by simp
next
    case False
    have norm-argument: norm (( $\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) z$ 
 $*_R y) + \text{indicat-real } (f - \{x\} \cap \text{space } M) z *_R x) = \text{norm } (\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) + \text{norm } (\text{indicat-real } (f - \{x\} \cap \text{space } M) z *_R x)$ 
    if  $z: z \in \text{space } M$  for  $z$ 
    proof (cases  $f z = x$ )
    case True
    have  $\text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y = 0$  if  $y \in F$  for  $y$  using
    True insert  $z$  that 1 unfolding indicator-def by force
    hence  $(\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) = 0$  by (meson
    sum.neutral)
    thus ?thesis by force
    qed (force)
    show ?thesis using False fin-0 fin-1 f norm-argument by (subst *, subst add,
    presburger add: insert, intro insert, intro insert, fastforce simp add: indicator-def
    intro!: insert(2) f(3), auto intro!: indicator insert nonneg-x)
    qed
next
    case 2
    show ?case
    proof (cases  $x = 0$ )
    case True
    then show ?thesis unfolding * using insert by simp
    next
    case False
    then show ?thesis unfolding * using insert simple-functionD(2)[OF f(1)]
by fast
    qed
next
    case 3
    show ?case
    proof (cases  $x = 0$ )
    case True
    then show ?thesis unfolding * using insert by simp
    next
    case False
    have  $\text{emeasure } M \{y \in \text{space } M. (\sum x \in \text{insert } x F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \leq \text{emeasure } M \{y \in \text{space } M. (\sum x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\}$ 
    using ** simple-functionD(2)[OF insert(6)] simple-functionD(2)[OF f(1)]
    insert.IH(2) by (intro emeasure-mono, blast, simp)
    also have  $\dots \leq \text{emeasure } M \{y \in \text{space } M. (\sum x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} + \text{emeasure } M \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\}$ 
    using simple-functionD(2)[OF insert(6)] simple-functionD(2)[OF f(1)]

```

by (*intro emeasure-subadditive, force+*)
 also have $\dots < \infty$ **using** *insert(7) fin-1[OF False]* **by** (*simp add: less-top*)
 finally show *?thesis* **by** *simp*
 qed
next
 case (ξ)
 thus *?case* **using** *insert nonneg-x f(3)* **by** (*auto simp add: scaleR-nonneg-nonneg*
intro: sum-nonneg)
 qed
moreover have *simple-function* $M (\lambda x. \sum_{y \in f^{-1} \text{space } M} \text{indicat-real } (f - \{y\})$
 $\cap \text{space } M) x *_R y)$ **using** *calculation* **by** *blast*
moreover have $P (\lambda x. \sum_{y \in f^{-1} \text{space } M} \text{indicat-real } (f - \{y\}) \cap \text{space } M) x$
 $*_R y)$ **using** *calculation* **by** *blast*
moreover have $\bigwedge x. x \in \text{space } M \implies 0 \leq f x$ **using** *f(3)* **by** *simp*
ultimately show *?thesis* **by** (*intro cong[OF - - - f(1,2)]*, *blast*, *blast*, *fast*)
presburger+
qed

lemma *finite-nn-integral-imp-ae-finite:*

fixes $f :: 'a \Rightarrow \text{ennreal}$
assumes $f \in \text{borel-measurable } M (\int^+ x. f x \partial M) < \infty$
shows $AE\ x\ \text{in } M. f x < \infty$
proof (*rule ccontr, goal-cases*)
case 1
let $?A = \text{space } M \cap \{x. f x = \infty\}$
have $*$: $\text{emeasure } M\ ?A > 0$ **using** 1 *assms(1)* **by** (*metis (mono-tags, lifting)*
assms(2) eventually-mono infinity-ennreal-def nn-integral-not-eq-infinite top.not-eq-extremum)
have $(\int^+ x. f x * \text{indicator } ?A\ x\ \partial M) = (\int^+ x. \infty * \text{indicator } ?A\ x\ \partial M)$ **by**
(metis (mono-tags, lifting) indicator-inter-arith indicator-simps(2) mem-Collect-eq
mult-eq-0-iff)
also have $\dots = \infty * \text{emeasure } M\ ?A$ **using** *assms(1)* **by** (*intro nn-integral-cmult-indicator,*
simp)
also have $\dots = \infty$ **using** $*$ **by** *fastforce*
finally have $(\int^+ x. f x * \text{indicator } ?A\ x\ \partial M) = \infty$.
moreover have $(\int^+ x. f x * \text{indicator } ?A\ x\ \partial M) \leq (\int^+ x. f x \partial M)$ **by** (*intro*
nn-integral-mono, simp add: indicator-def)
ultimately show *?case* **using** *assms(2)* **by** *simp*
qed

lemma *cauchy-L1-AE-cauchy-subseq:*

fixes $s :: \text{nat} \Rightarrow 'a \Rightarrow 'b :: \{\text{banach, second-countable-topology}\}$
assumes [*measurable*]: $\bigwedge n. \text{integrable } M\ (s\ n)$
and $\bigwedge e. e > 0 \implies \exists N. \forall i \geq N. \forall j \geq N. \text{LINT } x | M. \text{norm } (s\ i\ x - s\ j\ x) < e$
obtains r **where** *strict-mono* $r\ AE\ x\ \text{in } M. \text{Cauchy } (\lambda i. s\ (r\ i)\ x)$
proof–

have $\exists r. \forall n. (\forall i \geq r n. \forall j \geq r n. \text{LINT } x | M. \text{norm } (s \ i \ x - s \ j \ x) < (1 / 2) \wedge n) \wedge (r \ (\text{Suc } n) > r \ n)$
proof (intro dependent-nat-choice, goal-cases)
case 1
then show ?case **using** assms(2) **by** presburger
next
case (2 x n)
obtain N **where** *: $\text{LINT } x | M. \text{norm } (s \ i \ x - s \ j \ x) < (1 / 2) \wedge \text{Suc } n$ **if** $i \geq N \ j \geq N$ **for** i j **using** assms(2)[of (1 / 2) \wedge Suc n] **by** auto
{
fix i j **assume** $i \geq \max N \ (\text{Suc } x) \ j \geq \max N \ (\text{Suc } x)$
hence $\text{integral}^L M \ (\lambda x. \text{norm } (s \ i \ x - s \ j \ x)) < (1 / 2) \wedge \text{Suc } n$ **using** * **by** fastforce
}
then show ?case **by** fastforce
qed
then obtain r **where** strict-mono: strict-mono r **and** $\forall i \geq r n. \forall j \geq r n. \text{LINT } x | M. \text{norm } (s \ i \ x - s \ j \ x) < (1 / 2) \wedge n$ **for** n **using** strict-mono-Suc-iff **by** blast
hence r-is: $\text{LINT } x | M. \text{norm } (s \ (r \ (\text{Suc } n)) \ x - s \ (r \ n) \ x) < (1 / 2) \wedge n$ **for** n **by** (simp add: strict-mono-leD)

define g **where** $g = (\lambda n \ x. (\sum i \leq n. \text{ennreal } (\text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x))))$
define g' **where** $g' = (\lambda x. \sum i. \text{ennreal } (\text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x)))$

have integrable-g: $(\int^+ x. g \ n \ x \ \partial M) < 2$ **for** n
proof -
have $(\int^+ x. g \ n \ x \ \partial M) = (\int^+ x. (\sum i \leq n. \text{ennreal } (\text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x))) \ \partial M)$ **using** g-def **by** simp
also have ... = $(\sum i \leq n. (\int^+ x. \text{ennreal } (\text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x)) \ \partial M))$ **by** (intro nn-integral-sum, simp)
also have ... = $(\sum i \leq n. \text{LINT } x | M. \text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x))$
unfolding dist-norm **using** assms(1) **by** (subst nn-integral-eq-integral[OF integrable-norm], auto)
also have ... < $\text{ennreal } (\sum i \leq n. (1 / 2) \wedge i)$ **by** (intro ennreal-lessI[OF sum-pos sum-strict-mono[OF finite-atMost - r-is]], auto)
also have ... $\leq \text{ennreal } 2$ **unfolding** sum-gp0[$\text{of } 1 / 2 \ n$] **by** (intro ennreal-leI, auto)
finally show $(\int^+ x. g \ n \ x \ \partial M) < 2$ **by** simp
qed

have integrable-g': $(\int^+ x. g' \ x \ \partial M) \leq 2$
proof -
have incseq $(\lambda n. g \ n \ x)$ **for** x **by** (intro incseq-SucI, auto simp add: g-def ennreal-leI)
hence convergent $(\lambda n. g \ n \ x)$ **for** x **unfolding** convergent-def **using** LIMSEQ-incseq-SUP **by** fast
hence $(\lambda n. g \ n \ x) \longrightarrow g' \ x$ **for** x **unfolding** g-def g'-def **by** (intro summable-iff-convergent'[THEN iffD2, THEN summable-LIMSEQ], blast)

hence $(\int^+ x. g' x \partial M) = (\int^+ x. \liminf (\lambda n. g n x) \partial M)$ **by** (*metis lim-imp-Liminf trivial-limit-sequentially*)
also have $\dots \leq \liminf (\lambda n. \int^+ x. g n x \partial M)$ **by** (*intro nn-integral-liminf, simp add: g-def*)
also have $\dots \leq \liminf (\lambda n. 2)$ **using** *integrable-g* **by** (*intro Liminf-mono*) (*simp add: order-le-less*)
also have $\dots = 2$ **using** *sequentially-bot tendsto-iff-Liminf-eq-Limsup* **by** *blast*
finally show *?thesis* .
qed
hence *AE x in M. g' x < ∞* **by** (*intro finite-nn-integral-imp-ae-finite*) (*auto simp add: order-le-less-trans g'-def*)
moreover have *summable (λi. norm (s (r (Suc i)) x - s (r i) x))* **if** $g' x \neq \infty$ **for** x **using** *that unfolding g'-def* **by** (*intro summable-suminf-not-top*) *fastforce* +

ultimately have *ae-summable: AE x in M. summable (λi. s (r (Suc i)) x - s (r i) x)* **using** *summable-norm-cancel* **unfolding** *dist-norm* **by** *force*

{
 fix x **assume** *summable (λi. s (r (Suc i)) x - s (r i) x)*
 hence $(\lambda n. \sum i < n. s (r (Suc i)) x - s (r i) x) \longrightarrow (\sum i. s (r (Suc i)) x - s (r i) x)$ **using** *summable-LIMSEQ* **by** *blast*
 moreover have $(\lambda n. (\sum i < n. s (r (Suc i)) x - s (r i) x)) = (\lambda n. s (r n) x - s (r 0) x)$ **using** *sum-lessThan-telescope* **by** *fastforce*
 ultimately have $(\lambda n. s (r n) x - s (r 0) x) \longrightarrow (\sum i. s (r (Suc i)) x - s (r i) x)$ **by** *argo*
 hence $(\lambda n. s (r n) x - s (r 0) x + s (r 0) x) \longrightarrow (\sum i. s (r (Suc i)) x - s (r i) x) + s (r 0) x$ **by** (*intro isCont-tendsto-compose[of - λz. z + s (r 0) x], auto*)
 hence *Cauchy (λn. s (r n) x)* **by** (*simp add: LIMSEQ-imp-Cauchy*)
}
hence *AE x in M. Cauchy (λi. s (r i) x)* **using** *ae-summable* **by** *fast*
thus *?thesis* **by** (*rule that[OF strict-mono(1)]*)
qed

3.2 Linearly Ordered Banach Spaces

lemma *integrable-max[simp, intro]:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology}\}$

assumes *fg[measurable]: integrable M f integrable M g*

shows *integrable M (λx. max (f x) (g x))*

proof (*rule Bochner-Integration.integrable-bound*)

{
 fix $x y :: 'b$
 have $\text{norm } (\max x y) \leq \max (\text{norm } x) (\text{norm } y)$ **by** *linarith*
 also have $\dots \leq \text{norm } x + \text{norm } y$ **by** *simp*
 finally have $\text{norm } (\max x y) \leq \text{norm } (\text{norm } x + \text{norm } y)$ **by** *simp*
}
thus *AE x in M. norm (max (f x) (g x)) ≤ norm (norm (f x) + norm (g x))* **by** *simp*
qed (*auto intro: Bochner-Integration.integrable-add[OF integrable-norm[OF fg(1)]]*)

integrable-norm[*OF fg*(2)])])

lemma *integrable-min*[*simp*, *intro*]:

fixes *f* :: 'a \Rightarrow 'b :: {*second-countable-topology*, *banach*, *linorder-topology*}

assumes [*measurable*]: *integrable M f integrable M g*

shows *integrable M* ($\lambda x. \min (f x) (g x)$)

proof –

have *norm* ($\min (f x) (g x)$) \leq *norm* (*f x*) \vee *norm* ($\min (f x) (g x)$) \leq *norm* (*g x*) **for** *x* **by** *linarith*

thus ?thesis **by** (*intro integrable-bound*[*OF integrable-max*[*OF integrable-norm*(1,1), *OF assms*]], *auto*)

qed

lemma *integral-nonneg-AE-banach*:

fixes *f* :: 'a \Rightarrow 'b :: {*second-countable-topology*, *banach*, *linorder-topology*, *ordered-real-vector*}

assumes [*measurable*]: *f* \in *borel-measurable M* **and** *nonneg*: *AE x in M. 0 \leq f x*

shows $0 \leq \text{integral}^L M f$

proof *cases*

assume *integrable*: *integrable M f*

hence *max*: ($\lambda x. \max 0 (f x)$) \in *borel-measurable M* $\wedge x. 0 \leq \max 0 (f x)$ *integrable M* ($\lambda x. \max 0 (f x)$) **by** *auto*

hence $0 \leq \text{integral}^L M (\lambda x. \max 0 (f x))$

proof –

obtain *s* **where** *: $\wedge i. \text{simple-function } M (s i)$

$\wedge i. \text{emeasure } M \{y \in \text{space } M. s i y \neq 0\} \neq \infty$

$\wedge x. x \in \text{space } M \implies (\lambda i. s i x) \longrightarrow f x$

$\wedge x i. x \in \text{space } M \implies \text{norm } (s i x) \leq 2 * \text{norm } (f x)$ **using**

integrable-implies-simple-function-sequence[*OF integrable*] **by** *blast*

have *simple*: $\wedge i. \text{simple-function } M (\lambda x. \max 0 (s i x))$ **using** * **by** *fast*

have $\wedge i. \{y \in \text{space } M. \max 0 (s i y) \neq 0\} \subseteq \{y \in \text{space } M. s i y \neq 0\}$

unfolding *max-def* **by** *force*

moreover **have** $\wedge i. \{y \in \text{space } M. s i y \neq 0\} \in \text{sets } M$ **using** * **by** *measurable*

ultimately **have** $\wedge i. \text{emeasure } M \{y \in \text{space } M. \max 0 (s i y) \neq 0\} \leq \text{emeasure } M \{y \in \text{space } M. s i y \neq 0\}$ **by** (*rule emeasure-mono*)

hence **: $\wedge i. \text{emeasure } M \{y \in \text{space } M. \max 0 (s i y) \neq 0\} \neq \infty$ **using** *(2) **by** (*auto intro: order.strict-trans1 simp add: top.not-eq-extremum*)

have $\wedge x. x \in \text{space } M \implies (\lambda i. \max 0 (s i x)) \longrightarrow \max 0 (f x)$ **using** *(3) *tendsto-max* **by** *blast*

moreover **have** $\wedge x i. x \in \text{space } M \implies \text{norm } (\max 0 (s i x)) \leq \text{norm } (2 *_{\mathbb{R}} f x)$ **using** *(4) **unfolding** *max-def* **by** *auto*

ultimately **have** *tendsto*: $(\lambda i. \text{integral}^L M (\lambda x. \max 0 (s i x))) \longrightarrow \text{integral}^L M (\lambda x. \max 0 (f x))$

using *borel-measurable-simple-function simple integrable* **by** (*intro integral-dominated-convergence*[*OF max*(1) - *integrable-norm*[*OF integrable-scaleR-right*], *of - 2 f*], *blast+*)

{

fix *h* :: 'a \Rightarrow 'b :: {*second-countable-topology*, *banach*, *linorder-topology*, *ordered-real-vector*}

```

    assume simple-function M h emeasure M {y ∈ space M. h y ≠ 0} ≠ ∞ ∧ x.
x ∈ space M → h x ≥ 0
    hence *: integralL M h ≥ 0
    proof (induct rule: integrable-simple-function-induct-nn)
      case (cong f g)
      then show ?case using Bochner-Integration.integral-cong by force
    next
      case (indicator A y)
      hence A ≠ {} ⇒ y ≥ 0 using sets.sets-into-space by fastforce
      then show ?case using indicator by (cases A = {}, auto simp add:
scaleR-nonneg-nonneg)
    next
      case (add f g)
      then show ?case by (simp add: integrable-simple-function)
    qed
  }
  thus ?thesis using LIMSEQ-le-const[OF tendsto, of 0] ** simple by fastforce
qed
also have ... = integralL M f using nonneg by (auto intro: integral-cong-AE)
finally show ?thesis .
qed (simp add: not-integrable-integral-eq)

```

lemma *integral-mono-AE-banach*:

```

  fixes f g :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
dered-real-vector}
  assumes integrable M f integrable M g AE x in M. f x ≤ g x
  shows integralL M f ≤ integralL M g
  using integral-nonneg-AE-banach[of λx. g x - f x] assms Bochner-Integration.integral-diff[OF
assms(1,2)] by force

```

lemma *integral-mono-banach*:

```

  fixes f g :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
dered-real-vector}
  assumes integrable M f integrable M g ∧ x. x ∈ space M ⇒ f x ≤ g x
  shows integralL M f ≤ integralL M g
  using integral-mono-AE-banach assms by blast

```

3.3 Integrability and Measurability of the Diameter

context

```

  fixes s :: nat ⇒ 'a ⇒ 'b :: {second-countable-topology, banach} and M
  assumes bounded: ∧x. x ∈ space M ⇒ bounded (range (λi. s i x))
begin

```

lemma *borel-measurable-diameter*:

```

  assumes [measurable]: ∧i. (s i) ∈ borel-measurable M
  shows (λx. diameter {s i x | i. n ≤ i}) ∈ borel-measurable M
proof -
  have {s i x | i. N ≤ i} ≠ {} for x N by blast

```


hence *diameter-SUP*: $\text{diameter } \{s \ i \ x \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s \ i \ x) (s \ j \ x)) \text{ for } x \ N \text{ unfolding diameter-def by (auto intro!: arg-cong[of - Sup])}$

have *case-prod dist* ‘ $(\{s \ i \ x \mid i. n \leq i\} \times \{s \ i \ x \mid i. n \leq i\}) = ((\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) \text{ ‘ } (\{n..\} \times \{n..\}))$ ’ **for** x **by** *fast*

hence *: $(\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\}) = (\lambda x. \text{Sup } ((\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) \text{ ‘ } (\{n..\} \times \{n..\})))$ **using** *diameter-SUP* **by** (*simp add: case-prod-beta*)

have *bounded* $((\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) \text{ ‘ } (\{n..\} \times \{n..\}))$ **if** $x \in \text{space } M$ **for** x **by** (*rule bounded-imp-dist-bounded[OF bounded, OF that]*)

hence *bdd*: *bdd-above* $((\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) \text{ ‘ } (\{n..\} \times \{n..\}))$ **if** $x \in \text{space } M$ **for** x **using** *that bounded-imp-bdd-above* **by** *presburger*

have *fst* $p \in \text{borel-measurable } M$ *snd* $p \in \text{borel-measurable } M$ **if** $p \in s \text{ ‘ } \{n..\} \times s \text{ ‘ } \{n..\}$ **for** p **using** *that* **by** *fastforce+*

hence $(\lambda x. \text{fst } p \ x - \text{snd } p \ x) \in \text{borel-measurable } M$ **if** $p \in s \text{ ‘ } \{n..\} \times s \text{ ‘ } \{n..\}$ **for** p **using** *that borel-measurable-diff* **by** *simp*

hence $(\lambda x. \text{case } p \text{ of } (f, g) \Rightarrow \text{dist } (f \ x) (g \ x)) \in \text{borel-measurable } M$ **if** $p \in s \text{ ‘ } \{n..\} \times s \text{ ‘ } \{n..\}$ **for** p **unfolding** *dist-norm* **using** *that* **by** *measurable*

moreover **have** *countable* $(s \text{ ‘ } \{n..\} \times s \text{ ‘ } \{n..\})$ **by** (*intro countable-SIGMA countable-image, auto*)

ultimately show *?thesis* **unfolding** * **by** (*auto intro!: borel-measurable-cSUP bdd*)
qed

lemma *integrable-bound-diameter*:

fixes $f :: 'a \Rightarrow \text{real}$

assumes *integrable* $M \ f$

and [*measurable*]: $\bigwedge i. (s \ i) \in \text{borel-measurable } M$

and $\bigwedge x \ i. x \in \text{space } M \implies \text{norm } (s \ i \ x) \leq f \ x$

shows *integrable* $M \ (\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\})$

proof –

have $\{s \ i \ x \mid i. N \leq i\} \neq \{\}$ **for** $x \ N$ **by** *blast*

hence *diameter-SUP*: $\text{diameter } \{s \ i \ x \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s \ i \ x) (s \ j \ x)) \text{ for } x \ N \text{ unfolding diameter-def by (auto intro!: arg-cong[of - Sup])}$

{

fix x **assume** $x: x \in \text{space } M$

let $?S = (\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) \text{ ‘ } (\{n..\} \times \{n..\})$

have *case-prod dist* ‘ $(\{s \ i \ x \mid i. n \leq i\} \times \{s \ i \ x \mid i. n \leq i\}) = (\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) \text{ ‘ } (\{n..\} \times \{n..\})$ ’ **by** *fast*

hence *: $\text{diameter } \{s \ i \ x \mid i. n \leq i\} = \text{Sup } ?S$ **using** *diameter-SUP* **by** (*simp add: case-prod-beta*)

have *bounded* $?S$ **by** (*rule bounded-imp-dist-bounded[OF bounded[OF x]]*)

hence *Sup-S-nonneg:0* $\leq \text{Sup } ?S$ **by** (*auto intro!: cSup-upper2 x bounded-imp-bdd-above*)

have $\text{dist } (s \ i \ x) (s \ j \ x) \leq 2 * f \ x$ **for** $i \ j$ **by** (*intro dist-triangle2[THEN order-trans, of - 0]*) (*metis norm-conv-dist assms(3) x add-mono mult-2*)

```

    hence  $\forall c \in ?S. c \leq 2 * f x$  by force
    hence  $Sup ?S \leq 2 * f x$  by (intro cSup-least, auto)
    hence  $norm (Sup ?S) \leq 2 * norm (f x)$  using Sup-S-nonneg by auto
    also have  $\dots = norm (2 *_R f x)$  by simp
    finally have  $norm (diameter \{s\ i\ x \mid i. n \leq i\}) \leq norm (2 *_R f x)$  unfolding
* .
}
    hence  $\forall x \in M. norm (diameter \{s\ i\ x \mid i. n \leq i\}) \leq norm (2 *_R f x)$  by blast
    thus  $integrable\ M (\lambda x. diameter \{s\ i\ x \mid i. n \leq i\})$  using borel-measurable-diameter
by (intro Bochner-Integration.integrable-bound[OF assms(1)[THEN integrable-scaleR-right[of
2]]], measurable)
qed
end

end
theory Set-Integral-Addendum
  imports HOL-Analysis.Set-Integral Bochner-Integration-Addendum
begin

```

4 Auxiliary Lemmas for Integrals on a Set

```

lemma set-integral-scaleR-left:
  assumes  $A \in sets\ M\ c \neq 0 \implies integrable\ M\ f$ 
  shows  $LINT\ t:A|M. f\ t *_R c = (LINT\ t:A|M. f\ t) *_R c$ 
  unfolding set-lebesgue-integral-def
  using integrable-mult-indicator[OF assms]
  by (subst integral-scaleR-left[symmetric], auto)

lemma nn-set-integral-eq-set-integral:
  assumes [measurable]:  $integrable\ M\ f$ 
  and  $\forall x \in A\ in\ M. 0 \leq f\ x\ A \in sets\ M$ 
  shows  $(\int^+ x \in A. f\ x\ \partial M) = (\int x \in A. f\ x\ \partial M)$ 
proof -
  have  $(\int^+ x. indicator\ A\ x *_R f\ x\ \partial M) = (\int x \in A. f\ x\ \partial M)$ 
  unfolding set-lebesgue-integral-def using assms(2) by (intro nn-integral-eq-integral[of
-  $\lambda x. indicat\_real\ A\ x *_R f\ x$ ], blast intro: assms integrable-mult-indicator, fastforce)
  moreover have  $(\int^+ x. indicator\ A\ x *_R f\ x\ \partial M) = (\int^+ x \in A. f\ x\ \partial M)$  by (metis
ennreal-0 indicator-simps(1) indicator-simps(2) mult.commute mult-1 mult-zero-left
real-scaleR-def)
  ultimately show ?thesis by argo
qed

lemma set-integral-restrict-space:
  fixes  $f :: 'a \Rightarrow 'b::\{banach, second-countable-topology\}$ 
  assumes  $\Omega \cap space\ M \in sets\ M$ 
  shows  $set\_lebesgue\_integral\ (restrict\_space\ M\ \Omega)\ A\ f = set\_lebesgue\_integral\ M\ A$ 
 $(\lambda x. indicator\ \Omega\ x *_R f\ x)$ 
  unfolding set-lebesgue-integral-def
  by (subst integral-restrict-space, auto intro!: integrable-mult-indicator assms simp:

```

mult.commute)

lemma *set-integral-const*:

fixes $c :: 'b :: \{\text{banach}, \text{second-countable-topology}\}$
assumes $A \in \text{sets } M \text{ emeasure } M \ A \neq \infty$
shows $\text{set-lebesgue-integral } M \ A \ (\lambda \cdot. c) = \text{measure } M \ A *_R c$
unfolding *set-lebesgue-integral-def*
using *assms* **by** (*metis has-bochner-integral-indicator has-bochner-integral-integral-eq infinity-enreal-def less-top*)

lemma *set-integral-mono-banach*:

fixes $f \ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}, \text{linorder-topology}, \text{ordered-real-vector}\}$
assumes $\text{set-integrable } M \ A \ f \ \text{set-integrable } M \ A \ g$
 $\bigwedge x. x \in A \implies f \ x \leq g \ x$
shows $(\text{LINT } x:A | M. f \ x) \leq (\text{LINT } x:A | M. g \ x)$
using *assms* **unfolding** *set-integrable-def set-lebesgue-integral-def*
by (*auto intro: integral-mono-banach split: split-indicator*)

lemma *set-integral-mono-AE-banach*:

fixes $f \ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}, \text{linorder-topology}, \text{ordered-real-vector}\}$
assumes $\text{set-integrable } M \ A \ f \ \text{set-integrable } M \ A \ g \ \text{AE } x \in A \text{ in } M. f \ x \leq g \ x$
shows $\text{set-lebesgue-integral } M \ A \ f \leq \text{set-lebesgue-integral } M \ A \ g$ **using** *assms*
unfolding *set-lebesgue-integral-def* **by** (*auto simp add: set-integrable-def intro!: integral-mono-AE-banach[of M $\lambda x. \text{indicator } A \ x *_R f \ x \ \lambda x. \text{indicator } A \ x *_R g \ x]$, simp add: indicator-def*)

end

theory *Sigma-Finite-Measure-Addendum*

imports *Set-Integral-Addendum*

begin

5 Averaging Theorem

lemma *balls-countable-basis*:

obtains $D :: 'a :: \{\text{metric-space}, \text{second-countable-topology}\} \text{ set}$
where $\text{topological-basis } (\text{case-prod ball } ' (D \times (\mathbb{Q} \cap \{0 < ..\})))$
and $\text{countable } D$
and $D \neq \{\}$
proof –
obtain $D :: 'a \text{ set}$ **where** $\text{dense-subset: countable } D \ D \neq \{\} \llbracket \text{open } U; U \neq \{\} \rrbracket$
 $\implies \exists y \in D. y \in U$ **for** U **using** *countable-dense-exists* **by** *blast*
have $\text{topological-basis } (\text{case-prod ball } ' (D \times (\mathbb{Q} \cap \{0 < ..\})))$
proof (*intro topological-basis-iff[THEN iffD2], fast, clarify*)
fix U **and** $x :: 'a$ **assume** $\text{asm: open } U \ x \in U$
obtain e **where** $e: e > 0 \text{ ball } x \ e \subseteq U$ **using** *asm openE* **by** *blast*
obtain y **where** $y: y \in D \ y \in \text{ball } x \ (e / 3)$ **using** $\text{dense-subset}(3)[\text{OF open-ball, of } x \ e / 3]$ *centre-in-ball[THEN iffD2, OF divide-pos-pos[OF e(1), of 3]]* **by** *force*

obtain r **where** $r: r \in \mathbb{Q} \cap \{e/3 < \dots < e/2\}$ **unfolding** *Rats-def* **using** *of-rat-dense*[*OF*
divide-strict-left-mono[*OF* - $e(1)$], *of 2 3*] **by** *auto*

have $*$: $x \in \text{ball } y \ r$ **using** $r \ y$ **by** (*simp add: dist-commute*)
hence $\text{ball } y \ r \subseteq U$ **using** r **by** (*intro order-trans*[*OF* - $e(2)$], *simp*, *metric*)
moreover **have** $\text{ball } y \ r \in (\text{case-prod ball } ' (D \times (\mathbb{Q} \cap \{0 < \dots\})))$ **using** $y(1)$
 r **by** *force*
ultimately show $\exists B' \in (\text{case-prod ball } ' (D \times (\mathbb{Q} \cap \{0 < \dots\})))$. $x \in B' \wedge B' \subseteq$
 U **using** $*$ **by** *meson*
qed
thus *?thesis* **using** *that dense-subset* **by** *blast*
qed

context *sigma-finite-measure*
begin

lemma *sigma-finite-measure-induct*[*case-names finite-measure*, *consumes 0*]:

assumes $\bigwedge(N :: 'a \text{ measure}) \ \Omega. \text{finite-measure } N$
 $\implies N = \text{restrict-space } M \ \Omega$
 $\implies \Omega \in \text{sets } M$
 $\implies \text{emeasure } N \ \Omega \neq \infty$
 $\implies \text{emeasure } N \ \Omega \neq 0$
 $\implies \text{almost-everywhere } N \ Q$

and [*measurable*]: *Measurable.pred* $M \ Q$

shows *almost-everywhere* $M \ Q$

proof –

have $*$: *almost-everywhere* $N \ Q$ **if** *finite-measure* $N \ N = \text{restrict-space } M \ \Omega \ \Omega$
 $\in \text{sets } M \ \text{emeasure } N \ \Omega \neq \infty$ **for** $N \ \Omega$ **using** *that* **by** (*cases* *emeasure* $N \ \Omega = 0$,
auto *intro: emeasure-0-AE assms(1)*)

obtain $A :: \text{nat} \Rightarrow 'a \text{ set}$ **where** $A: \text{range } A \subseteq \text{sets } M \ (\bigcup i. A \ i) = \text{space } M$ **and**
emeasure-finite: emeasure $M \ (A \ i) \neq \infty$ **for** i **using** *sigma-finite* **by** *metis*

note $A(1)$ [*measurable*]

have *space-restr: space* (*restrict-space* $M \ (A \ i)$) = $A \ i$ **for** i **unfolding** *space-restrict-space*
by *simp*

{

fix i

have $*$: $\{x \in A \ i \cap \text{space } M. \ Q \ x\} = \{x \in \text{space } M. \ Q \ x\} \cap (A \ i)$ **by** *fast*

have *Measurable.pred* (*restrict-space* $M \ (A \ i)$) Q **using** A **by** (*intro measurableI*,
auto *simp add: space-restr intro!: sets-restrict-space-iff*[*THEN iffD2*], *measurable*,
auto)

}

note *this*[*measurable*]

{

fix i

have *finite-measure* (*restrict-space* $M \ (A \ i)$) **using** *emeasure-finite* **by** (*intro*
finite-measureI, *subst space-restr*, *subst emeasure-restrict-space*, *auto*)

hence *emeasure* (*restrict-space* $M \ (A \ i)$) $\{x \in A \ i. \neg Q \ x\} = 0$ **using** *emea-*
sure-finite **by** (*intro AE-iff-measurable*[*THEN iffD1*, *OF* - $*$], *measurable*, *subst*

$\text{space-restr}[\text{symmetric}], \text{intro sets.top}, \text{auto simp add: emeasure-restrict-space})$
hence $\text{emeasure } M \{x \in A \mid i. \neg Q x\} = 0$ **by** ($\text{subst emeasure-restrict-space}[\text{symmetric}],$
 $\text{auto})$
 $\}$
hence $\text{emeasure } M (\bigcup i. \{x \in A \mid i. \neg Q x\}) = 0$ **by** ($\text{intro emeasure-UN-eq-0},$
 $\text{auto})$
moreover have $(\bigcup i. \{x \in A \mid i. \neg Q x\}) = \{x \in \text{space } M. \neg Q x\}$ **using** A **by**
 auto
ultimately show $?thesis$ **by** ($\text{intro AE-iff-measurable}[\text{THEN iffD2}], \text{auto})$
qed

lemma averaging-theorem:

fixes $f::\Rightarrow 'b::\{\text{second-countable-topology}, \text{banach}\}$
assumes $[\text{measurable}]: \text{integrable } M f$
and $\text{closed: closed } S$
and $\bigwedge A. A \in \text{sets } M \implies \text{measure } M A > 0 \implies (1 / \text{measure } M A) *_{\mathbb{R}}$
 $\text{set-lebesgue-integral } M A f \in S$
shows $\text{AE } x \text{ in } M. f x \in S$
proof ($\text{induct rule: sigma-finite-measure-induct})$
case ($\text{finite-measure } N \Omega$)

interpret $\text{finite-measure } N$ **by** ($\text{rule finite-measure}$)

have $\text{integrable}[\text{measurable}]: \text{integrable } N f$ **using** $\text{assms finite-measure}$ **by** (auto
 $\text{simp: integrable-restrict-space integrable-mult-indicator})$
have $\text{average: } (1 / \text{Sigma-Algebra.measure } N A) *_{\mathbb{R}} \text{set-lebesgue-integral } N A f$
 $\in S$ **if** $A \in \text{sets } N$ $\text{measure } N A > 0$ **for** A
proof –
have $*$: $A \in \text{sets } M$ **using** that **by** ($\text{simp add: sets-restrict-space-iff finite-measure}$)
have $A = A \cap \Omega$ **by** ($\text{metis finite-measure}(2,3) \text{ inf.orderE sets.sets-into-space}$
 $\text{space-restrict-space that}(1))$
hence $\text{set-lebesgue-integral } N A f = \text{set-lebesgue-integral } M A f$ **unfolding**
 finite-measure **by** ($\text{subst set-integral-restrict-space}, \text{auto simp add: finite-measure}$
 $\text{set-lebesgue-integral-def indicator-inter-arith}[\text{symmetric}])$
moreover have $\text{measure } N A = \text{measure } M A$ **using** that **by** ($\text{auto intro!:$
 $\text{measure-restrict-space simp add: finite-measure sets-restrict-space-iff})$
ultimately show $?thesis$ **using** $\text{that } * \text{ assms}(3)$ **by** presburger
qed

obtain $D :: 'b \text{ set}$ **where** $\text{balls-basis: topological-basis } (\text{case-prod ball } ' (D \times (\mathbb{Q} \cap \{0 < ..\})))$ **and** $\text{countable-D: countable } D$ **using** $\text{balls-countable-basis}$ **by** blast
have $\text{countable-balls: countable } (\text{case-prod ball } ' (D \times (\mathbb{Q} \cap \{0 < ..\})))$ **using**
 $\text{countable-rat countable-D}$ **by** blast

obtain B **where** $B\text{-balls: } B \subseteq \text{case-prod ball } ' (D \times (\mathbb{Q} \cap \{0 < ..\})) \cup B = -S$
using $\text{topological-basis}[\text{THEN iffD1}, \text{OF balls-basis}] \text{ open-Compl}[\text{OF assms}(2)]$ **by**
 meson

hence $\text{countable-B: countable } B$ **using** $\text{countable-balls countable-subset}$ **by** fast

```

define  $b$  where  $b = \text{from-nat-into } (B \cup \{\{\}\})$ 
have  $B \cup \{\{\}\} \neq \{\}$  by simp
have  $\text{range-}b$ :  $\text{range } b = B \cup \{\{\}\}$  using countable-B by (auto simp add: b-def
intro!: range-from-nat-into)
have  $\text{open-}b$ :  $\text{open } (b \ i)$  for  $i$  unfolding  $b\text{-def}$  using B-balls open-ball from-nat-into[of
 $B \cup \{\{\}\} \ i$ ] by force
have  $\text{Union-range-}b$ :  $\bigcup (\text{range } b) = -S$  using B-balls range-b by simp

{
  fix  $v \ r$  assume ball-in-Compl:  $\text{ball } v \ r \subseteq -S$ 
  define  $A$  where  $A = f^{-1} \text{ball } v \ r \cap \text{space } N$ 
  have  $\text{dist-less}$ :  $\text{dist } (f \ x) \ v < r$  if  $x \in A$  for  $x$  using that unfolding  $A\text{-def}$ 
vimage-def by (simp add: dist-commute)
  hence  $AE\text{-less}$ :  $AE \ x \in A \text{ in } N. \text{norm } (f \ x - v) < r$  by (auto simp add:
dist-norm)
  have  $*$ :  $A \in \text{sets } N$  unfolding  $A\text{-def}$  by simp
  have  $\text{emeasure } N \ A = 0$ 
  proof -
  {
    assume asm:  $\text{emeasure } N \ A > 0$ 
    hence  $\text{measure-pos}$ :  $\text{measure } N \ A > 0$  unfolding emeasure-eq-measure by
simp
    hence  $(1 / \text{measure } N \ A) *_R \text{set-lebesgue-integral } N \ A \ f - v = (1 / \text{measure } N$ 
 $A) *_R \text{set-lebesgue-integral } N \ A \ (\lambda x. f \ x - v)$  using integrable integrable-const * by
(subst set-integral-diff(2), auto simp add: set-integrable-def set-integral-const[OF *]
algebra-simps intro!: integrable-mult-indicator)
    moreover have  $\text{norm } (\int_{x \in A}. (f \ x - v) \partial N) \leq (\int_{x \in A}. \text{norm } (f \ x$ 
 $- v) \partial N)$  using  $*$  by (auto intro!: integral-norm-bound[of  $N \ \lambda x. \text{indicator } A \ x$ 
 $*_R (f \ x - v)$ , THEN order-trans] integrable-mult-indicator integrable simp add:
set-lebesgue-integral-def)
    ultimately have  $\text{norm } ((1 / \text{measure } N \ A) *_R \text{set-lebesgue-integral } N \ A \ f$ 
 $- v) \leq \text{set-lebesgue-integral } N \ A \ (\lambda x. \text{norm } (f \ x - v)) / \text{measure } N \ A$  using asm
by (auto intro: divide-right-mono)
    also have  $\dots < \text{set-lebesgue-integral } N \ A \ (\lambda x. r) / \text{measure } N \ A$ 
    unfolding set-lebesgue-integral-def
    using asm * integrable integrable-const AE-less measure-pos
by (intro divide-strict-right-mono integral-less-AE[of  $- \ A$ ] integrable-mult-indicator)
    (fastforce simp add: dist-less dist-norm indicator-def) +
    also have  $\dots = r$  using  $\text{measure-pos}$  by (simp add: set-integral-const)
    finally have  $\text{dist } ((1 / \text{measure } N \ A) *_R \text{set-lebesgue-integral } N \ A \ f) \ v < r$ 
by (subst dist-norm)
    hence False using average[OF * measure-pos] by (metis ComplD dist-commute
in-mono mem-ball ball-in-Compl)
  }
  thus ?thesis by fastforce
qed
}
note  $*$  = this

```

```

{
  fix b' assume b' ∈ B
  hence ball-subset-Compl: b' ⊆ -S and ball-radius-pos: ∃ v ∈ D. ∃ r > 0. b' =
ball v r using B-balls by (blast, fast)
}
note ** = this
hence emeasure N (f -' b i ∩ space N) = 0 for i by (cases b i = {}, simp)
(metis UnE singletonD * range-b[THEN eq-refl, THEN range-subsetD])
hence emeasure N (⋃ i. f -' b i ∩ space N) = 0 using open-b by (intro
emeasure-UN-eq-0) fastforce+
moreover have (⋃ i. f -' b i ∩ space N) = f -' (⋃ (range b)) ∩ space N by
blast
ultimately have emeasure N (f -' (-S) ∩ space N) = 0 using Union-range-b
by argo
hence AE x in N. f x ∉ -S using open-Compl[OF assms(2)] by (intro AE-iff-measurable[THEN
iffD2], auto)
thus ?case by force
qed (simp add: pred-sets2[OF borel-closed] assms(2))

```

lemma density-zero:

```

fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach}
assumes integrable M f
and density-0: ⋀ A. A ∈ sets M ⇒ set-lebesgue-integral M A f = 0
shows AE x in M. f x = 0
using averaging-theorem[OF assms(1), of {0}] assms(2)
by (simp add: scaleR-nonneg-nonneg)

```

lemma density-unique:

```

fixes f f' :: 'a ⇒ 'b :: {second-countable-topology, banach}
assumes integrable M f integrable M f'
and density-eq: ⋀ A. A ∈ sets M ⇒ set-lebesgue-integral M A f = set-lebesgue-integral
M A f'
shows AE x in M. f x = f' x

```

proof—

```

{
  fix A assume asm: A ∈ sets M
  hence LINT x | M. indicat-real A x *R (f x - f' x) = 0 using density-eq
assms(1,2) by (simp add: set-lebesgue-integral-def algebra-simps Bochner-Integration.integral-diff[OF
integrable-mult-indicator(1,1)])
}
thus ?thesis using density-zero[OF Bochner-Integration.integrable-diff[OF assms(1,2)]]
by (simp add: set-lebesgue-integral-def)
qed

```

lemma density-nonneg:

```

fixes f :: - ⇒ 'b :: {second-countable-topology, banach, linorder-topology, ordered-real-vector}
assumes integrable M f
and ⋀ A. A ∈ sets M ⇒ set-lebesgue-integral M A f ≥ 0
shows AE x in M. f x ≥ 0

```

```

using averaging-theorem[OF assms(1), of {0..}, OF closed-atLeast] assms(2)
by (simp add: scaleR-nonneg-nonneg)

corollary integral-nonneg-AE-eq-0-iff-AE:
  fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$ 
  assumes  $f[\text{measurable}]: \text{integrable } M f$  and  $\text{nonneg}: AE\ x\ in\ M. 0 \leq f\ x$ 
  shows  $\text{integral}^L\ M\ f = 0 \iff (AE\ x\ in\ M. f\ x = 0)$ 
proof
  assume *:  $\text{integral}^L\ M\ f = 0$ 
  {
    fix  $A$  assume  $asm: A \in \text{sets } M$ 
    have  $0 \leq \text{integral}^L\ M\ (\lambda x. \text{indicator } A\ x *_{\mathbb{R}} f\ x)$  using  $\text{nonneg}$  by (subst integral-zero[of  $M$ , symmetric], intro integral-mono-AE-banach integrable-mult-indicator  $asm\ f$  integrable-zero, auto simp add: indicator-def)
    moreover have  $\dots \leq \text{integral}^L\ M\ f$  using  $\text{nonneg}$  by (intro integral-mono-AE-banach integrable-mult-indicator  $asm\ f$ , auto simp add: indicator-def)
    ultimately have  $\text{set-lebesgue-integral } M\ A\ f = 0$  unfolding set-lebesgue-integral-def using * by force
  }
  thus  $AE\ x\ in\ M. f\ x = 0$  by (intro density-zero  $f$ , blast)
qed (auto simp add: integral-eq-zero-AE)

corollary integral-eq-mono-AE-eq-AE:
  fixes  $f\ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$ 
  assumes  $\text{integrable } M\ f\ \text{integrable } M\ g$   $\text{integral}^L\ M\ f = \text{integral}^L\ M\ g$   $AE\ x\ in\ M. f\ x \leq g\ x$ 
  shows  $AE\ x\ in\ M. f\ x = g\ x$ 
proof –
  define  $h$  where  $h = (\lambda x. g\ x - f\ x)$ 
  have  $AE\ x\ in\ M. h\ x = 0$  unfolding  $h\text{-def}$  using  $assms$  by (subst integral-nonneg-AE-eq-0-iff-AE[symmetric]) auto
  then show ?thesis unfolding  $h\text{-def}$  by auto
qed

end

end

theory Conditional-Expectation-Banach
imports HOL-Probability.Conditional-Expectation Sigma-Finite-Measure-Addendum
begin

```

6 Conditional Expectation in Banach Spaces

```

definition has-cond-exp ::  $'a\ \text{measure} \Rightarrow 'a\ \text{measure} \Rightarrow ('a \Rightarrow 'b) \Rightarrow ('a \Rightarrow 'b :: \{\text{real-normed-vector, second-countable-topology}\}) \Rightarrow \text{bool}$  where
   $\text{has-cond-exp } M\ F\ f\ g = ((\forall A \in \text{sets } F. (\int x \in A. f\ x\ \partial M) = (\int x \in A. g\ x\ \partial M)))$ 

```


$\wedge \text{integrable } M f$
 $\wedge \text{integrable } M g$
 $\wedge g \in \text{borel-measurable } F$)

lemma *has-cond-expI'*:
assumes $\bigwedge A. A \in \text{sets } F \implies (\int x \in A. f x \partial M) = (\int x \in A. g x \partial M)$
 $\text{integrable } M f$
 $\text{integrable } M g$
 $g \in \text{borel-measurable } F$
shows *has-cond-exp* $M F f g$
using *assms* **unfolding** *has-cond-exp-def* **by** *simp*

lemma *has-cond-expD*:
assumes *has-cond-exp* $M F f g$
shows $\bigwedge A. A \in \text{sets } F \implies (\int x \in A. f x \partial M) = (\int x \in A. g x \partial M)$
 $\text{integrable } M f$
 $\text{integrable } M g$
 $g \in \text{borel-measurable } F$
using *assms* **unfolding** *has-cond-exp-def* **by** *simp+*

definition *cond-exp* :: $'a \text{ measure} \Rightarrow 'a \text{ measure} \Rightarrow ('a \Rightarrow 'b) \Rightarrow ('a \Rightarrow 'b :: \{\text{banach}, \text{second-countable-topology}\})$ **where**
 $\text{cond-exp } M F f = (\text{if } \exists g. \text{has-cond-exp } M F f g \text{ then } (\text{SOME } g. \text{has-cond-exp } M F f g) \text{ else } (\lambda -. 0))$

lemma *borel-measurable-cond-exp[measurable]*: $\text{cond-exp } M F f \in \text{borel-measurable } F$
by (*metis cond-exp-def someI has-cond-exp-def borel-measurable-const*)

lemma *integrable-cond-exp[intro]*: $\text{integrable } M (\text{cond-exp } M F f)$
by (*metis cond-exp-def has-cond-expD(3) integrable-zero someI*)

lemma *set-integrable-cond-exp[intro]*:
assumes $A \in \text{sets } M$
shows *set-integrable* $M A (\text{cond-exp } M F f)$ **using** *integrable-mult-indicator[OF assms integrable-cond-exp, of F f]* **by** (*auto simp add: set-integrable-def intro!: integrable-mult-indicator[OF assms integrable-cond-exp]*)

context *sigma-finite-subalgebra*
begin

lemma *borel-measurable-cond-exp'[measurable]*: $\text{cond-exp } M F f \in \text{borel-measurable } M$
by (*metis cond-exp-def someI has-cond-exp-def borel-measurable-const subalg measurable-from-subalg*)

lemma *cond-exp-null*:

assumes $\nexists g. \text{has-cond-exp } M F f g$
shows $\text{cond-exp } M F f = (\lambda \cdot. 0)$
unfolding cond-exp-def **using** assms **by** argo

lemma $\text{has-cond-exp-nested-subalg}$:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes $\text{subalgebra } G F \text{ has-cond-exp } M F f h \text{ has-cond-exp } M G f h'$
shows $\text{has-cond-exp } M F h' h$
by $(\text{intro has-cond-expI'}) (\text{metis assms has-cond-expD in-mono subalgebra-def})+$

lemma $\text{has-cond-exp-charact}$:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes $\text{has-cond-exp } M F f g$
shows $\text{has-cond-exp } M F f (\text{cond-exp } M F f)$
 $AE x \text{ in } M. \text{cond-exp } M F f x = g x$

proof –

show $\text{cond-exp: has-cond-exp } M F f (\text{cond-exp } M F f)$ **using** assms someI
 cond-exp-def **by** metis

let $?MF = \text{restr-to-subalg } M F$

interpret $\text{sigma-finite-measure } ?MF$ **by** $(\text{rule sigma-fin-subalg})$

{

fix A **assume** $A \in \text{sets } ?MF$

then have $[\text{measurable}]: A \in \text{sets } F$ **using** $\text{sets-restr-to-subalg}[OF \text{ subalg}]$ **by**

simp

have $(\int x \in A. g x \partial ?MF) = (\int x \in A. g x \partial M)$ **using** assms subalg **by** $(\text{auto simp add: integral-subalgebra2 set-lebesgue-integral-def dest!: has-cond-expD})$

also have $\dots = (\int x \in A. \text{cond-exp } M F f x \partial M)$ **using** assms cond-exp **by** $(\text{simp add: has-cond-exp-def})$

also have $\dots = (\int x \in A. \text{cond-exp } M F f x \partial ?MF)$ **using** subalg **by** $(\text{auto simp add: integral-subalgebra2 set-lebesgue-integral-def})$

finally have $(\int x \in A. g x \partial ?MF) = (\int x \in A. \text{cond-exp } M F f x \partial ?MF)$ **by** simp

}

hence $AE x \text{ in } ?MF. \text{cond-exp } M F f x = g x$ **using** $\text{cond-exp assms subalg}$ **by** $(\text{intro density-unique, auto dest: has-cond-expD intro!: integrable-in-subalg})$

then show $AE x \text{ in } M. \text{cond-exp } M F f x = g x$ **using** $AE\text{-restr-to-subalg}[OF \text{ subalg}]$ **by** simp

qed

lemma cond-exp-charact :

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$

assumes $\bigwedge A. A \in \text{sets } F \implies (\int x \in A. f x \partial M) = (\int x \in A. g x \partial M)$

$\text{integrable } M f$

$\text{integrable } M g$

$g \in \text{borel-measurable } F$

shows $AE x \text{ in } M. \text{cond-exp } M F f x = g x$

by $(\text{intro has-cond-exp-charact has-cond-expI' assms}) \text{ auto}$

corollary $\text{cond-exp-F-meas}[\text{intro, simp}]$:

fixes $f :: 'a \Rightarrow 'b::\{\text{second-countable-topology, banach}\}$
assumes $\text{integrable } M f$
 $f \in \text{borel-measurable } F$
shows $AE\ x\ \text{in } M. \text{cond-exp } M\ F\ f\ x = f\ x$
by (rule cond-exp-charact , auto intro: assms)

Congruence

lemma has-cond-exp-cong :
assumes $\text{integrable } M f \wedge x. x \in \text{space } M \implies f\ x = g\ x$ $\text{has-cond-exp } M\ F\ g\ h$
shows $\text{has-cond-exp } M\ F\ f\ h$
proof (intro has-cond-expI' [OF - $\text{assms}(1)$], goal-cases)
case (1 A)
hence $\text{set-lebesgue-integral } M\ A\ f = \text{set-lebesgue-integral } M\ A\ g$ **by** (intro $\text{set-lebesgue-integral-cong}$)
(meson $\text{assms}(2)$ subalg in-mono subalgebra-def sets.sets-into-space subalgebra-def subsetD)+
then show ?case **using** 1 $\text{assms}(3)$ **by** (simp add: has-cond-exp-def)
qed (auto simp add: has-cond-expD [OF $\text{assms}(3)$])

lemma cond-exp-cong :
fixes $f :: 'a \Rightarrow 'b::\{\text{second-countable-topology, banach}\}$
assumes $\text{integrable } M f \text{ integrable } M g \wedge x. x \in \text{space } M \implies f\ x = g\ x$
shows $AE\ x\ \text{in } M. \text{cond-exp } M\ F\ f\ x = \text{cond-exp } M\ F\ g\ x$
proof (cases $\exists h. \text{has-cond-exp } M\ F\ f\ h$)
case True
then obtain h **where** $h: \text{has-cond-exp } M\ F\ f\ h \text{ has-cond-exp } M\ F\ g\ h$ **using**
 $\text{has-cond-exp-cong assms}$ **by** metis
show ?thesis **using** $h[\text{THEN } \text{has-cond-exp-charact}(2)]$ **by** fastforce
next
case False
moreover have $\nexists h. \text{has-cond-exp } M\ F\ g\ h$ **using** False $\text{has-cond-exp-cong assms}$
by auto
ultimately show ?thesis **unfolding** cond-exp-def **by** auto
qed

lemma $\text{has-cond-exp-cong-AE}$:
assumes $\text{integrable } M f\ AE\ x\ \text{in } M. f\ x = g\ x$ $\text{has-cond-exp } M\ F\ g\ h$
shows $\text{has-cond-exp } M\ F\ f\ h$
using $\text{assms}(1,2)$ subalg subalgebra-def subset-iff
by (intro has-cond-expI' , subst $\text{set-lebesgue-integral-cong-AE}$ [OF - $\text{assms}(1)$ [THEN
 $\text{borel-measurable-integrable}$] $\text{borel-measurable-integrable}(1)$ [OF $\text{has-cond-expD}(2)$ [OF
 $\text{assms}(3)$]]])
(fast intro: has-cond-expD [OF $\text{assms}(3)$] $\text{integrable-cong-AE-imp}$ [OF - - AE-symmetric]) +

lemma $\text{has-cond-exp-cong-AE'}$:
assumes $h \in \text{borel-measurable } F\ AE\ x\ \text{in } M. h\ x = h'\ x$ $\text{has-cond-exp } M\ F\ f\ h'$
shows $\text{has-cond-exp } M\ F\ f\ h$
using $\text{assms}(1, 2)$ subalg subalgebra-def subset-iff
using $\text{AE-restr-to-subalg2}$ [OF subalg $\text{assms}(2)$] measurable-from-subalg
by (intro has-cond-expI' , subst $\text{set-lebesgue-integral-cong-AE}$ [OF - measurable-from-subalg(1,1)] [OF

subalg], *OF* - *assms*(1) *has-cond-expD*(4)[*OF* *assms*(3)]]
 (*fast intro*: *has-cond-expD*[*OF* *assms*(3)] *integrable-cong-AE-imp*[*OF* - - *AE-symmetric*])+

lemma *cond-exp-cong-AE*:
 fixes *f* :: 'a \Rightarrow 'b::{*second-countable-topology*,*banach*}
 assumes *integrable* *M* *f* *integrable* *M* *g* *AE* *x* in *M*. *f* *x* = *g* *x*
 shows *AE* *x* in *M*. *cond-exp* *M* *F* *f* *x* = *cond-exp* *M* *F* *g* *x*
proof (*cases* \exists *h*. *has-cond-exp* *M* *F* *f* *h*)
 case *True*
 then obtain *h* where *h*: *has-cond-exp* *M* *F* *f* *h* *has-cond-exp* *M* *F* *g* *h* **using**
has-cond-exp-cong-AE *assms* **by** (*metis* (*mono-tags*, *lifting*) *eventually-mono*)
 show ?thesis **using** *h*[*THEN* *has-cond-exp-charact*(2)] **by** *fastforce*
 next
 case *False*
 moreover have \nexists *h*. *has-cond-exp* *M* *F* *g* *h* **using** *False* *has-cond-exp-cong-AE*
assms **by** *auto*
 ultimately show ?thesis **unfolding** *cond-exp-def* **by** *auto*
qed

lemma *has-cond-exp-real*:
 fixes *f* :: 'a \Rightarrow *real*
 assumes *integrable* *M* *f*
 shows *has-cond-exp* *M* *F* *f* (*real-cond-exp* *M* *F* *f*)
by (*intro* *has-cond-expI'*, *auto* *intro!*: *real-cond-exp-intA* *assms*)

lemma *cond-exp-real*[*intro*]:
 fixes *f* :: 'a \Rightarrow *real*
 assumes *integrable* *M* *f*
 shows *AE* *x* in *M*. *cond-exp* *M* *F* *f* *x* = *real-cond-exp* *M* *F* *f* *x*
using *has-cond-exp-charact* *has-cond-exp-real* *assms* **by** *blast*

lemma *cond-exp-cmult*:
 fixes *f* :: 'a \Rightarrow *real*
 assumes *integrable* *M* *f*
 shows *AE* *x* in *M*. *cond-exp* *M* *F* ($\lambda x. c * f x$) *x* = *c* * *cond-exp* *M* *F* *f* *x*
using *real-cond-exp-cmult*[*OF* *assms*(1), *of* *c*] *assms*(1)[*THEN* *cond-exp-real*]
assms(1)[*THEN* *integrable-mult-right*, *THEN* *cond-exp-real*, *of* *c*] **by** *fastforce*

Indicator functions

lemma *has-cond-exp-indicator*:
 assumes *A* \in *sets* *M* *emeasure* *M* *A* < ∞
 shows *has-cond-exp* *M* *F* ($\lambda x. \text{indicat-real } A \ x *_{\mathbb{R}} y$) ($\lambda x. \text{real-cond-exp } M \ F$
 (*indicator* *A*) *x* * _{\mathbb{R}} *y*)
proof (*intro* *has-cond-expI'*, *goal-cases*)
 case (1 *B*)
 have $\int x \in B. (\text{indicat-real } A \ x *_{\mathbb{R}} y) \ \partial M = (\int x \in B. \text{indicat-real } A \ x \ \partial M) *_{\mathbb{R}}$
 y **using** *assms* **by** (*intro* *set-integral-scaleR-left*, *meson* 1 *in-mono* *subalg* *subalge-*
bra-def, *blast*)
 also have ... = $(\int x \in B. \text{real-cond-exp } M \ F (\text{indicator } A) \ x \ \partial M) *_{\mathbb{R}} y$ **using** 1

```

assms by (subst real-cond-exp-intA, auto)
  also have ... =  $\int_{x \in B}. (\text{real-cond-exp } M \ F \ (\text{indicator } A) \ x \ *_R \ y) \ \partial M$  using
assms by (intro set-integral-scaleR-left[symmetric], meson 1 in-mono subalg subal-
gebra-def, blast)
  finally show ?case .
next
  case 2
  then show ?case using integrable-scaleR-left integrable-real-indicator assms by
blast
next
  case 3
  show ?case using assms by (intro integrable-scaleR-left, intro real-cond-exp-int,
blast+)
next
  case 4
  then show ?case by (intro borel-measurable-scaleR, intro Conditional-Expectation.borel-measurable-cond-exp,
simp)
qed

```

```

lemma cond-exp-indicator[intro]:
  fixes  $y :: 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
  assumes [measurable]:  $A \in \text{sets } M$   $\text{emeasure } M \ A < \infty$ 
  shows  $\text{AE } x \text{ in } M. \text{cond-exp } M \ F \ (\lambda x. \text{indicat-real } A \ x \ *_R \ y) \ x = \text{cond-exp } M \ F$ 
 $(\text{indicator } A) \ x \ *_R \ y$ 
proof -
  have  $\text{AE } x \text{ in } M. \text{cond-exp } M \ F \ (\lambda x. \text{indicat-real } A \ x \ *_R \ y) \ x = \text{real-cond-exp } M \ F$ 
 $(\text{indicator } A) \ x \ *_R \ y$  using has-cond-exp-indicator[OF assms] has-cond-exp-charact
by blast
  thus ?thesis using cond-exp-real[OF integrable-real-indicator, OF assms] by fast-
force
qed

```

Addition

```

lemma has-cond-exp-add:
  fixes  $f \ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
  assumes has-cond-exp  $M \ F \ f \ f'$  has-cond-exp  $M \ F \ g \ g'$ 
  shows has-cond-exp  $M \ F \ (\lambda x. f \ x + g \ x) \ (\lambda x. f' \ x + g' \ x)$ 
proof (intro has-cond-expI', goal-cases)
  case (1 A)
  have  $\int_{x \in A}. (f \ x + g \ x) \partial M = (\int_{x \in A}. f \ x \ \partial M) + (\int_{x \in A}. g \ x \ \partial M)$  using
assms[THEN has-cond-expD(2)] subalg 1 by (intro set-integral-add(2), auto simp
add: subalgebra-def set-integrable-def intro: integrable-mult-indicator)
  also have ... =  $(\int_{x \in A}. f' \ x \ \partial M) + (\int_{x \in A}. g' \ x \ \partial M)$  using assms[THEN
has-cond-expD(1)[OF - 1]] by argo
  also have ... =  $\int_{x \in A}. (f' \ x + g' \ x) \partial M$  using assms[THEN has-cond-expD(3)]
subalg 1 by (intro set-integral-add(2)[symmetric], auto simp add: subalgebra-def
set-integrable-def intro: integrable-mult-indicator)
  finally show ?case .
next

```

```

    case 2
  then show ?case by (metis Bochner-Integration.integrable-add assms has-cond-expD(2))
next
  case 3
  then show ?case by (metis Bochner-Integration.integrable-add assms has-cond-expD(3))
next
  case 4
  then show ?case using assms borel-measurable-add has-cond-expD(4) by blast
qed

```

```

lemma has-cond-exp-scaleR-right:
  fixes f :: 'a  $\Rightarrow$  'b::{second-countable-topology,banach}
  assumes has-cond-exp M F f f'
  shows has-cond-exp M F ( $\lambda x. c *_R f x$ ) ( $\lambda x. c *_R f' x$ )
  using has-cond-expD[OF assms] by (intro has-cond-expI', auto)

```

```

lemma cond-exp-scaleR-right:
  fixes f :: 'a  $\Rightarrow$  'b::{second-countable-topology,banach}
  assumes integrable M f
  shows AE x in M. cond-exp M F ( $\lambda x. c *_R f x$ ) x = c *_R cond-exp M F f x
proof (cases  $\exists f'. \text{has-cond-exp M F f f'}$ )
  case True
  then show ?thesis using assms has-cond-exp-charact has-cond-exp-scaleR-right
  by metis
next
  case False
  show ?thesis
  proof (cases c = 0)
    case True
    then show ?thesis by simp
  next
    case c-nonzero: False
    have  $\nexists f'. \text{has-cond-exp M F } (\lambda x. c *_R f x) f'$ 
    proof (standard, goal-cases)
      case 1
      then obtain f' where f': has-cond-exp M F ( $\lambda x. c *_R f x$ ) f' by blast
      have has-cond-exp M F f ( $\lambda x. \text{inverse } c *_R f' x$ ) using has-cond-expD[OF
f'] divideR-right[OF c-nonzero] assms by (intro has-cond-expI', auto)
      then show ?case using False by blast
    qed
    then show ?thesis using cond-exp-null[OF False] cond-exp-null by force
  qed
qed

```

```

lemma cond-exp-uminus:
  fixes f :: 'a  $\Rightarrow$  'b::{second-countable-topology,banach}
  assumes integrable M f
  shows AE x in M. cond-exp M F ( $\lambda x. - f x$ ) x = - cond-exp M F f x
  using cond-exp-scaleR-right[OF assms, of -1] by force

```

corollary *has-cond-exp-simple:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *simple-function* $M f$ *emeasure* $M \{y \in \text{space } M. f y \neq 0\} \neq \infty$
shows *has-cond-exp* $M F f$ (*cond-exp* $M F f$)
using *assms*
proof (*induction rule: integrable-simple-function-induct*)
case (*cong* $f g$)
then show ?*case* **using** *has-cond-exp-cong* **by** (*metis* (*no-types*, *opaque-lifting*)
Bochner-Integration.integrable-cong *has-cond-expD*(2) *has-cond-exp-charact*(1))
next
case (*indicator* $A y$)
then show ?*case* **using** *has-cond-exp-charact*[*OF* *has-cond-exp-indicator*] **by** *fast*
next
case (*add* $u v$)
then show ?*case* **using** *has-cond-exp-add* *has-cond-exp-charact*(1) **by** *blast*
qed

lemma *cond-exp-contraction-real:*

fixes $f :: 'a \Rightarrow \text{real}$
assumes *integrable*[*measurable*]: *integrable* $M f$
shows $\text{AE } x \text{ in } M. \text{norm} (\text{cond-exp } M F f x) \leq \text{cond-exp } M F (\lambda x. \text{norm} (f x)) x$
proof–
have *int*: *integrable* $M (\lambda x. \text{norm} (f x))$ **using** *assms* **by** *blast*
have *: $\text{AE } x \text{ in } M. 0 \leq \text{cond-exp } M F (\lambda x. \text{norm} (f x)) x$ **using** *cond-exp-real*[*THEN* *AE-symmetric*, *OF* *integrable-norm*[*OF* *integrable*]] *real-cond-exp-ge-c*[*OF* *integrable-norm*[*OF* *integrable*], *of 0*] *norm-ge-zero* **by** *fastforce*
have **: $A \in \text{sets } F \implies \int x \in A. |f x| \partial M = \int x \in A. \text{real-cond-exp } M F (\lambda x. \text{norm} (f x)) x \partial M$ **for** A **unfolding** *real-norm-def* **using** *assms* *integrable-abs* *real-cond-exp-intA* **by** *blast*

have *norm-int*: $A \in \text{sets } F \implies (\int x \in A. |f x| \partial M) = (\int^+ x \in A. |f x| \partial M)$ **for** A
using *assms* **by** (*intro nn-set-integral-eq-set-integral*[*symmetric*], *blast*, *fastforce*)
(*meson subalg subalgebra-def subsetD*)

have $\text{AE } x \text{ in } M. \text{real-cond-exp } M F (\lambda x. \text{norm} (f x)) x \geq 0$ **using** *int real-cond-exp-ge-c*
by *force*

hence *cond-exp-norm-int*: $A \in \text{sets } F \implies (\int x \in A. \text{real-cond-exp } M F (\lambda x. \text{norm} (f x)) x \partial M) = (\int^+ x \in A. \text{real-cond-exp } M F (\lambda x. \text{norm} (f x)) x \partial M)$ **for** A **using** *assms* **by** (*intro nn-set-integral-eq-set-integral*[*symmetric*], *blast*, *fastforce*)
(*meson subalg subalgebra-def subsetD*)

have $A \in \text{sets } F \implies \int^+ x \in A. |f x| \partial M = \int^+ x \in A. \text{real-cond-exp } M F (\lambda x. \text{norm} (f x)) x \partial M$ **for** A **using** ** *norm-int cond-exp-norm-int* **by** (*auto simp* *add: nn-integral-set-ennreal*)

moreover **have** $(\lambda x. \text{ennreal } |f x|) \in \text{borel-measurable } M$ **by** *measurable*
moreover **have** $(\lambda x. \text{ennreal} (\text{real-cond-exp } M F (\lambda x. \text{norm} (f x)) x)) \in \text{borel-measurable } F$ **by** *measurable*

ultimately **have** $\text{AE } x \text{ in } M. \text{nn-cond-exp } M F (\lambda x. \text{ennreal } |f x|) x = \text{real-cond-exp}$

$M F (\lambda x. \text{norm } (f x)) x$ **by** (intro nn-cond-exp-charact[THEN AE-symmetric], auto)
hence $AE x \text{ in } M. \text{nn-cond-exp } M F (\lambda x. \text{ennreal } |f x|) x \leq \text{cond-exp } M F (\lambda x. \text{norm } (f x)) x$ **using** cond-exp-real[OF int] **by** force
moreover have $AE x \text{ in } M. |\text{real-cond-exp } M F f x| = \text{norm } (\text{cond-exp } M F f x)$
unfolding real-norm-def **using** cond-exp-real[OF assms] * **by** force
ultimately have $AE x \text{ in } M. \text{ennreal } (\text{norm } (\text{cond-exp } M F f x)) \leq \text{cond-exp } M F (\lambda x. \text{norm } (f x)) x$ **using** real-cond-exp-abs[OF assms[THEN borel-measurable-integrable]] **by** fastforce
hence $AE x \text{ in } M. \text{enn2real } (\text{ennreal } (\text{norm } (\text{cond-exp } M F f x))) \leq \text{enn2real } (\text{cond-exp } M F (\lambda x. \text{norm } (f x)) x)$ **using** ennreal-le-iff2 **by** force
thus ?thesis **using** * **by** fastforce
qed

lemma cond-exp-contraction-simple:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes simple-function $M f$ emeasure $M \{y \in \text{space } M. f y \neq 0\} \neq \infty$
shows $AE x \text{ in } M. \text{norm } (\text{cond-exp } M F f x) \leq \text{cond-exp } M F (\lambda x. \text{norm } (f x)) x$
using assms
proof (induction rule: integrable-simple-function-induct)
case (cong f g)
hence $ae: AE x \text{ in } M. f x = g x$ **by** blast
hence $AE x \text{ in } M. \text{cond-exp } M F f x = \text{cond-exp } M F g x$ **using** cong has-cond-exp-simple
by (subst cond-exp-cong-AE) (auto intro!: has-cond-expD(2))
hence $AE x \text{ in } M. \text{norm } (\text{cond-exp } M F f x) = \text{norm } (\text{cond-exp } M F g x)$ **by** force
moreover have $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. \text{norm } (f x)) x = \text{cond-exp } M F (\lambda x. \text{norm } (g x)) x$ **using** ae cong has-cond-exp-simple **by** (subst cond-exp-cong-AE) (auto dest: has-cond-expD)
ultimately show ?case **using** cong(6) **by** fastforce
next
case (indicator A y)
hence $AE x \text{ in } M. \text{cond-exp } M F (\lambda a. \text{indicator } A a *_R y) x = \text{cond-exp } M F (\text{indicator } A) x *_R y$ **by** blast
hence *: $AE x \text{ in } M. \text{norm } (\text{cond-exp } M F (\lambda a. \text{indicat-real } A a *_R y) x) \leq \text{norm } y * \text{cond-exp } M F (\lambda x. \text{norm } (\text{indicat-real } A x)) x$ **using** cond-exp-contraction-real[OF integrable-real-indicator, OF indicator] **by** fastforce

have $AE x \text{ in } M. \text{norm } y * \text{cond-exp } M F (\lambda x. \text{norm } (\text{indicat-real } A x)) x = \text{norm } y * \text{real-cond-exp } M F (\lambda x. \text{norm } (\text{indicat-real } A x)) x$ **using** cond-exp-real[OF integrable-real-indicator, OF indicator] **by** fastforce

moreover have $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. \text{norm } y * \text{norm } (\text{indicat-real } A x)) x = \text{real-cond-exp } M F (\lambda x. \text{norm } y * \text{norm } (\text{indicat-real } A x)) x$ **using** indicator **by** (intro cond-exp-real, auto)

ultimately have $AE x \text{ in } M. \text{norm } y * \text{cond-exp } M F (\lambda x. \text{norm } (\text{indicat-real } A x)) x = \text{cond-exp } M F (\lambda x. \text{norm } y * \text{norm } (\text{indicat-real } A x)) x$ **using** real-cond-exp-cmult[of $\lambda x. \text{norm } (\text{indicat-real } A x) \text{ norm } y$] indicator **by** fastforce

moreover have $(\lambda x. \text{norm } y * \text{norm } (\text{indicat-real } A x)) = (\lambda x. \text{norm } (\text{indicat-real } A x *_R y))$ **by** force

ultimately show ?case using * by force
next
case (add u v)
have AE x in M. norm (cond-exp M F (λa. u a + v a) x) = norm (cond-exp M F u x + cond-exp M F v x) using has-cond-exp-charact(2)[OF has-cond-exp-add, OF has-cond-exp-simple(1,1), OF add(1,2,3,4)] by fastforce
moreover have AE x in M. norm (cond-exp M F u x + cond-exp M F v x) ≤ norm (cond-exp M F u x) + norm (cond-exp M F v x) using norm-triangle-ineq by blast
moreover have AE x in M. norm (cond-exp M F u x) + norm (cond-exp M F v x) ≤ cond-exp M F (λx. norm (u x)) x + cond-exp M F (λx. norm (v x)) x using add(6,7) by fastforce
moreover have AE x in M. cond-exp M F (λx. norm (u x)) x + cond-exp M F (λx. norm (v x)) x = cond-exp M F (λx. norm (u x) + norm (v x)) x using integrable-simple-function[OF add(1,2)] integrable-simple-function[OF add(3,4)] by (intro has-cond-exp-charact(2)[OF has-cond-exp-add[OF has-cond-exp-charact(1,1)], THEN AE-symmetric], auto intro: has-cond-exp-real)
moreover have AE x in M. cond-exp M F (λx. norm (u x) + norm (v x)) x = cond-exp M F (λx. norm (u x + v x)) x using add(5) integrable-simple-function[OF add(1,2)] integrable-simple-function[OF add(3,4)] by (intro cond-exp-cong, auto)
ultimately show ?case by force
qed

lemma has-cond-exp-simple-lim:

fixes f :: 'a ⇒ 'b::{second-countable-topology, banach}
assumes integrable[measurable]: integrable M f
and ∧i. simple-function M (s i)
and ∧i. emeasure M {y ∈ space M. s i y ≠ 0} ≠ ∞
and ∧x. x ∈ space M ⇒ (λi. s i x) ⟶ f x
and ∧x i. x ∈ space M ⇒ norm (s i x) ≤ 2 * norm (f x)
obtains r
where has-cond-exp M F f (λx. lim (λi. cond-exp M F (s (r i)) x))
AE x in M. convergent (λi. cond-exp M F (s (r i)) x)
strict-mono r

proof –

have [measurable]: (s i) ∈ borel-measurable M for i using assms(2) by (simp add: borel-measurable-simple-function)
have integrable-s: integrable M (λx. s i x) for i using assms(2) assms(3) integrable-simple-function by blast
have integrable-4f: integrable M (λx. 4 * norm (f x)) using assms(1) by simp
have integrable-2f: integrable M (λx. 2 * norm (f x)) using assms(1) by simp
have integrable-2-cond-exp-norm-f: integrable M (λx. 2 * cond-exp M F (λx. norm (f x)) x) by fast

have emeasure M {y ∈ space M. s i y - s j y ≠ 0} ≤ emeasure M {y ∈ space M. s i y ≠ 0} + emeasure M {y ∈ space M. s j y ≠ 0} for i j using simple-functionD(2)[OF assms(2)] by (intro order-trans[OF emeasure-mono emeasure-subadditive], auto)

hence fin-sup: emeasure M {y ∈ space M. s i y - s j y ≠ 0} ≠ ∞ for

$i\ j$ **using** *assms(3)* **by** (*metis (mono-tags) ennreal-add-eq-top linorder-not-less top.not-eq-extremum infinity-ennreal-def*)

have *emeasure* $M \{y \in \text{space } M. \text{norm } (s\ i\ y - s\ j\ y) \neq 0\} \leq \text{emeasure } M \{y \in \text{space } M. s\ i\ y \neq 0\} + \text{emeasure } M \{y \in \text{space } M. s\ j\ y \neq 0\}$ **for** $i\ j$ **using** *simple-functionD(2)[OF assms(2)]* **by** (*intro order-trans[OF emeasure-mono emeasure-subadditive], auto*)

hence *fin-sup-norm*: *emeasure* $M \{y \in \text{space } M. \text{norm } (s\ i\ y - s\ j\ y) \neq 0\} \neq \infty$ **for** $i\ j$ **using** *assms(3)* **by** (*metis (mono-tags) ennreal-add-eq-top linorder-not-less top.not-eq-extremum infinity-ennreal-def*)

have *Cauchy*: *Cauchy* $(\lambda n. s\ n\ x)$ **if** $x \in \text{space } M$ **for** x **using** *assms(4)* *LIM-SEQ-imp-Cauchy* **that** **by** *blast*

hence *bounded-range-s*: *bounded* $(\text{range } (\lambda n. s\ n\ x))$ **if** $x \in \text{space } M$ **for** x **using** *that cauchy-imp-bounded* **by** *fast*

have *AE* x **in** $M. (\lambda n. \text{diameter } \{s\ i\ x \mid i. n \leq i\}) \longrightarrow 0$ **using** *Cauchy cauchy-iff-diameter-tends-to-zero-and-bounded* **by** *fast*

moreover **have** $(\lambda x. \text{diameter } \{s\ i\ x \mid i. n \leq i\}) \in \text{borel-measurable } M$ **for** n **using** *bounded-range-s borel-measurable-diameter* **by** *measurable*

moreover **have** *AE* x **in** $M. \text{norm } (\text{diameter } \{s\ i\ x \mid i. n \leq i\}) \leq 4 * \text{norm } (f\ x)$ **for** n

proof -

{

fix x **assume** $x: x \in \text{space } M$

have $\text{diameter } \{s\ i\ x \mid i. n \leq i\} \leq 2 * \text{norm } (f\ x) + 2 * \text{norm } (f\ x)$ **by** (*intro diameter-le, blast, subst dist-norm[symmetric], intro dist-triangle3[THEN order-trans, of 0], intro add-mono*) (*auto intro: assms(5)[OF x]*)

hence $\text{norm } (\text{diameter } \{s\ i\ x \mid i. n \leq i\}) \leq 4 * \text{norm } (f\ x)$ **using** *diameter-ge-0[OF bounded-subset[OF bounded-range-s], OF x, of {s i x | i. n ≤ i}]* **by** *force*

}

thus *?thesis* **by** *fast*

qed

ultimately **have** *diameter-tendsto-zero*: $(\lambda n. \text{LINT } x | M. \text{diameter } \{s\ i\ x \mid i. n \leq i\}) \longrightarrow 0$ **by** (*intro integral-dominated-convergence[OF borel-measurable-const[of 0] - integrable-4f, simplified]*) (*fast+*)

have *diameter-integrable*: *integrable* $M (\lambda x. \text{diameter } \{s\ i\ x \mid i. n \leq i\})$ **for** n **using** *assms(1,5)* **by** (*intro integrable-bound-diameter[OF bounded-range-s integrable-2f], auto*)

have *dist-integrable*: *integrable* $M (\lambda x. \text{dist } (s\ i\ x) (s\ j\ x))$ **for** $i\ j$

using *assms(5) dist-triangle3[of s i - - 0, THEN order-trans, OF add-mono, of - 2 * norm (f -)]*

by (*intro Bochner-Integration.integrable-bound[OF integrable-4f]*) *fastforce+*

hence *dist-norm-integrable*: *integrable* $M (\lambda x. \text{norm } (s\ i\ x - s\ j\ x))$ **for** $i\ j$ **unfolding** *dist-norm* **by** *presburger*

have $\exists N. \forall i \geq N. \forall j \geq N. \text{LINT } x|M. \text{norm } (\text{cond-exp } M F (s i) x - \text{cond-exp } M F (s j) x) < e$ **if** $e\text{-pos}: e > 0$ **for** e
proof –
obtain N **where** $*$: $\text{LINT } x|M. \text{diameter } \{s i x \mid i. n \leq i\} < e$ **if** $n \geq N$ **for** n **using** *that order-tendsto-iff* [*THEN iffD1, OF diameter-tendsto-zero, unfolded eventually-sequentially*] $e\text{-pos}$ **by** *presburger*
 $\{$
fix $i j x$ **assume** $\text{asm}: i \geq N j \geq N x \in \text{space } M$
have $\text{case-prod dist } '(\{s i x \mid i. N \leq i\} \times \{s i x \mid i. N \leq i\}) = \text{case-prod } (\lambda i j. \text{dist } (s i x) (s j x)) '(\{N..\} \times \{N..\})$ **by** *fast*
hence $\text{diameter } \{s i x \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s i x) (s j x))$ **unfolding** *diameter-def* **by** *auto*
moreover **have** $(\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s i x) (s j x)) \geq \text{dist } (s i x) (s j x)$ **using** *asm bounded-imp-bdd-above* [*OF bounded-imp-dist-bounded, OF bounded-range-s*] **by** (*intro cSup-upper, auto*)
ultimately **have** $\text{diameter } \{s i x \mid i. N \leq i\} \geq \text{dist } (s i x) (s j x)$ **by** *presburger*
 $\}$
hence $\text{LINT } x|M. \text{dist } (s i x) (s j x) < e$ **if** $i \geq N j \geq N$ **for** $i j$ **using** *that* $*$ **by** (*intro integral-mono* [*OF dist-integrable diameter-integrable, THEN order.strict-trans1*], *blast+*)
moreover **have** $\text{LINT } x|M. \text{norm } (\text{cond-exp } M F (s i) x - \text{cond-exp } M F (s j) x) \leq \text{LINT } x|M. \text{dist } (s i x) (s j x)$ **for** $i j$
proof –
have $\text{LINT } x|M. \text{norm } (\text{cond-exp } M F (s i) x - \text{cond-exp } M F (s j) x) = \text{LINT } x|M. \text{norm } (\text{cond-exp } M F (s i) x + - 1 *_R \text{cond-exp } M F (s j) x)$ **unfolding** *dist-norm* **by** *simp*
also **have** $\dots = \text{LINT } x|M. \text{norm } (\text{cond-exp } M F (\lambda x. s i x - s j x) x)$ **using** *has-cond-exp-charact*(2) [*OF has-cond-exp-add* [*OF - has-cond-exp-scaleR-right, OF has-cond-exp-charact*(1,1), *OF has-cond-exp-simple*(1,1) [*OF assms*(2,3)]]], *THEN AE-symmetric, of i - 1 j*] **by** (*intro integral-cong-AE*) *force+*
also **have** $\dots \leq \text{LINT } x|M. \text{cond-exp } M F (\lambda x. \text{norm } (s i x - s j x)) x$ **using** *cond-exp-contraction-simple* [*OF - fin-sup, of i j*] *integrable-cond-exp assms*(2) **by** (*intro integral-mono-AE, fast+*)
also **have** $\dots = \text{LINT } x|M. \text{norm } (s i x - s j x)$ **unfolding** *set-integral-space*(1) [*OF integrable-cond-exp, symmetric*] *set-integral-space* [*OF dist-norm-integrable, symmetric*] **by** (*intro has-cond-expD*(1) [*OF has-cond-exp-simple* [*OF - fin-sup-norm*], *symmetric*]) (*metis assms*(2) *simple-function-compose1 simple-function-diff, metis sets.top subalg subalgebra-def*)
finally **show** *?thesis* **unfolding** *dist-norm* .
qed
ultimately **show** *?thesis* **using** *order.strict-trans1* **by** *meson*
qed
then **obtain** r **where** *strict-mono-r: strict-mono r* **and** *AE-Cauchy: AE x in M. Cauchy* ($\lambda i. \text{cond-exp } M F (s (r i)) x$) **by** (*rule cauchy-L1-AE-cauchy-subseq* [*OF integrable-cond-exp*], *auto*)
hence *ae-lim-cond-exp: AE x in M. $(\lambda n. \text{cond-exp } M F (s (r n)) x) \longrightarrow \lim (\lambda n. \text{cond-exp } M F (s (r n)) x)$* **using** *Cauchy-convergent-iff convergent-LIMSEQ-iff*

by *fastforce*

have *cond-exp-bounded*: $AE\ x\ in\ M.\ norm\ (cond-exp\ M\ F\ (s\ (r\ n))\ x) \leq cond-exp\ M\ F\ (\lambda x.\ 2 * norm\ (f\ x))\ x\ \text{for}\ n$

proof –

have $AE\ x\ in\ M.\ norm\ (cond-exp\ M\ F\ (s\ (r\ n))\ x) \leq cond-exp\ M\ F\ (\lambda x.\ norm\ (s\ (r\ n)\ x))\ x\ \text{by}\ (rule\ cond-exp-contraction-simple\ [OF\ assms(2,3)])$

moreover **have** $AE\ x\ in\ M.\ real-cond-exp\ M\ F\ (\lambda x.\ norm\ (s\ (r\ n)\ x))\ x \leq real-cond-exp\ M\ F\ (\lambda x.\ 2 * norm\ (f\ x))\ x\ \text{using}\ integrable-s\ integrable-2f\ assms(5)\ \text{by}\ (intro\ real-cond-exp-mono,\ auto)$

ultimately **show** *?thesis* **using** *cond-exp-real*[*OF integrable-norm, OF integrable-s, of r n*] *cond-exp-real*[*OF integrable-2f*] **by** *force*

qed

have *lim-integrable*: $integrable\ M\ (\lambda x.\ lim\ (\lambda i.\ cond-exp\ M\ F\ (s\ (r\ i))\ x))\ \text{by}\ (intro\ integrable-dominated-convergence\ [OF\ -\ borel-measurable-cond-exp'\ integrable-cond-exp\ ae-lim-cond-exp\ cond-exp-bounded],\ simp)$

{

fix *A* **assume** *A-in-sets-F*: $A \in sets\ F$

have $AE\ x\ in\ M.\ norm\ (indicator\ A\ x *_{R}\ cond-exp\ M\ F\ (s\ (r\ n))\ x) \leq cond-exp\ M\ F\ (\lambda x.\ 2 * norm\ (f\ x))\ x\ \text{for}\ n$

proof –

have $AE\ x\ in\ M.\ norm\ (indicator\ A\ x *_{R}\ cond-exp\ M\ F\ (s\ (r\ n))\ x) \leq norm\ (cond-exp\ M\ F\ (s\ (r\ n))\ x)\ \text{unfolding}\ indicator-def\ \text{by}\ simp$

thus *?thesis* **using** *cond-exp-bounded*[*of n*] **by** *force*

qed

hence *lim-cond-exp-int*: $(\lambda n.\ LINT\ x:A|M.\ cond-exp\ M\ F\ (s\ (r\ n))\ x) \longrightarrow LINT\ x:A|M.\ lim\ (\lambda n.\ cond-exp\ M\ F\ (s\ (r\ n))\ x)$

using *ae-lim-cond-exp measurable-from-subalg*[*OF subalg borel-measurable-indicator, OF A-in-sets-F*] *cond-exp-bounded*

unfolding *set-lebesgue-integral-def*

by (*intro integral-dominated-convergence*[*OF borel-measurable-scaleR borel-measurable-scaleR integrable-cond-exp*]) (*fastforce simp add: tendsto-scaleR*)+

have $AE\ x\ in\ M.\ norm\ (indicator\ A\ x *_{R}\ s\ (r\ n)\ x) \leq 2 * norm\ (f\ x)\ \text{for}\ n$

proof –

have $AE\ x\ in\ M.\ norm\ (indicator\ A\ x *_{R}\ s\ (r\ n)\ x) \leq norm\ (s\ (r\ n)\ x)$

unfolding *indicator-def* **by** *simp*

thus *?thesis* **using** *assms(5)*[*of - r n*] **by** *fastforce*

qed

hence *lim-s-int*: $(\lambda n.\ LINT\ x:A|M.\ s\ (r\ n)\ x) \longrightarrow LINT\ x:A|M.\ f\ x$

using *measurable-from-subalg*[*OF subalg borel-measurable-indicator, OF A-in-sets-F*] *LIMSEQ-subseq-LIMSEQ*[*OF assms(4) strict-mono-r*] *assms(5)*

unfolding *set-lebesgue-integral-def comp-def*

by (*intro integral-dominated-convergence*[*OF borel-measurable-scaleR borel-measurable-scaleR integrable-2f*]) (*fastforce simp add: tendsto-scaleR*)+

have $LINT\ x:A|M.\ lim\ (\lambda n.\ cond-exp\ M\ F\ (s\ (r\ n))\ x) = lim\ (\lambda n.\ LINT\ x:A|M.\ cond-exp\ M\ F\ (s\ (r\ n))\ x)\ \text{using}\ limI\ [OF\ lim-cond-exp-int]\ \text{by}\ argo$

also have ... = $\lim (\lambda n. \text{LINT } x:A | M. s (r n) x)$ **using** *has-cond-expD(1)* [*OF has-cond-exp-simple* [*OF assms(2,3)*] *A-in-sets-F*, *symmetric*] **by** *presburger*
also have ... = $\text{LINT } x:A | M. f x$ **using** *limI* [*OF lim-s-int*] **by** *argo*
finally have $\text{LINT } x:A | M. \lim (\lambda n. \text{cond-exp } M F (s (r n)) x) = \text{LINT } x:A | M. f x$.
}
hence *has-cond-exp* $M F f (\lambda x. \lim (\lambda i. \text{cond-exp } M F (s (r i)) x))$ **using** *assms(1)* *lim-integrable* **by** (*intro has-cond-expI'*, *auto*)
thus thesis using *AE-Cauchy Cauchy-convergent strict-mono-r* **by** (*auto intro!*: *that*)
qed

lemma *cond-exp-simple-lim*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes [*measurable*]: *integrable* $M f$
and $\bigwedge i. \text{simple-function } M (s i)$
and $\bigwedge i. \text{emeasure } M \{y \in \text{space } M. s i y \neq 0\} \neq \infty$
and $\bigwedge x. x \in \text{space } M \implies (\lambda i. s i x) \longrightarrow f x$
and $\bigwedge x i. x \in \text{space } M \implies \text{norm } (s i x) \leq 2 * \text{norm } (f x)$
obtains r **where** $\text{AE } x \text{ in } M. (\lambda i. \text{cond-exp } M F (s (r i)) x) \longrightarrow \text{cond-exp } M F f x$ *strict-mono r*
proof –
obtain r **where** $\text{AE } x \text{ in } M. \text{cond-exp } M F f x = \lim (\lambda i. \text{cond-exp } M F (s (r i)) x)$ $\text{AE } x \text{ in } M. \text{convergent } (\lambda i. \text{cond-exp } M F (s (r i)) x)$ *strict-mono r* **using** *has-cond-exp-charact(2)* **by** (*auto intro*: *has-cond-exp-simple-lim* [*OF assms*])
thus *?thesis* **by** (*auto intro!*: *that* [*of r*] *simp: convergent-LIMSEQ-iff*)
qed

corollary *has-cond-expI*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes *integrable* $M f$
shows *has-cond-exp* $M F f (\text{cond-exp } M F f)$
proof –
obtain s **where** $s\text{-is}: \bigwedge i. \text{simple-function } M (s i) \bigwedge i. \text{emeasure } M \{y \in \text{space } M. s i y \neq 0\} \neq \infty \bigwedge x. x \in \text{space } M \implies (\lambda i. s i x) \longrightarrow f x \bigwedge x i. x \in \text{space } M \implies \text{norm } (s i x) \leq 2 * \text{norm } (f x)$ **using** *integrable-implies-simple-function-sequence* [*OF assms*] **by** *blast*
show *?thesis* **using** *has-cond-exp-simple-lim* [*OF assms s-is*] *has-cond-exp-charact(1)*
by *metis*
qed

lemma *cond-exp-nested-subalg*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes *integrable* $M f$ *subalgebra* $M G$ *subalgebra* $G F$
shows $\text{AE } \xi \text{ in } M. \text{cond-exp } M F f \xi = \text{cond-exp } M F (\text{cond-exp } M G f) \xi$
using *has-cond-expI* *assms sigma-finite-subalgebra-def* **by** (*auto intro!*: *has-cond-exp-nested-subalg* [*THEN has-cond-exp-charact(2)*, *THEN AE-symmetric*] *sigma-finite-subalgebra.has-cond-expI* [*OF*

sigma-finite-subalgebra.intro[*OF assms*(2)] *nested-subalg-is-sigma-finite*)

lemma *cond-exp-set-integral*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f A \in \text{sets } F$
shows $(\int x \in A. f x \partial M) = (\int x \in A. \text{cond-exp } M F f x \partial M)$
using *has-cond-expD*(1)[*OF has-cond-expI*, *OF assms*] **by** *argo*

lemma *cond-exp-add*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f$ *integrable* $M g$
shows $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x + g x) x = \text{cond-exp } M F f x + \text{cond-exp } M F g x$
using *has-cond-exp-add*[*OF has-cond-expI*(1,1), *OF assms*, *THEN has-cond-exp-charact*(2)]
.

lemma *cond-exp-diff*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f$ *integrable* $M g$
shows $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x - g x) x = \text{cond-exp } M F f x - \text{cond-exp } M F g x$
using *has-cond-exp-add*[*OF - has-cond-exp-scaleR-right*, *OF has-cond-expI*(1,1), *OF assms*, *THEN has-cond-exp-charact*(2), *of -1*] **by** *simp*

lemma *cond-exp-diff'*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f$ *integrable* $M g$
shows $AE x \text{ in } M. \text{cond-exp } M F (f - g) x = \text{cond-exp } M F f x - \text{cond-exp } M F g x$
unfolding *fun-diff-def* **using** *assms* **by** (*rule cond-exp-diff*)

lemma *cond-exp-scaleR-left*:

fixes $f :: 'a \Rightarrow \text{real}$
assumes *integrable* $M f$
shows $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x *_R c) x = \text{cond-exp } M F f x *_R c$
using *cond-exp-set-integral*[*OF assms*] *subalg assms* **unfolding** *subalgebra-def*
by (*intro cond-exp-charact*,
subst set-integral-scaleR-left, *blast*, *intro assms*,
subst set-integral-scaleR-left, *blast*, *intro integrable-cond-exp*)
auto

lemma *cond-exp-contraction*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f$
shows $AE x \text{ in } M. \text{norm } (\text{cond-exp } M F f x) \leq \text{cond-exp } M F (\lambda x. \text{norm } (f x))$
x

proof –

obtain s **where** $s: \bigwedge i. \text{simple-function } M (s i) \bigwedge i. \text{emeasure } M \{y \in \text{space } M. s i y \neq 0\} \neq \infty \bigwedge x. x \in \text{space } M \implies (\lambda i. s i x) \longrightarrow f x \bigwedge i. x \in \text{space } M$

$\Rightarrow \text{norm } (s \ i \ x) \leq 2 * \text{norm } (f \ x)$

by (*blast intro: integrable-implies-simple-function-sequence*[*OF assms*])

obtain *r* **where** *r*: *AE x in M. (λi. cond-exp M F (s (r i)) x) ⟶ cond-exp M F f x strict-mono r* **using** *cond-exp-simple-lim*[*OF assms s*] **by** *blast*

have *norm-s-r*: $\bigwedge i. \text{simple-function } M \ (\lambda x. \text{norm } (s \ (r \ i) \ x)) \ \bigwedge i. \text{emeasure } M \ \{y \in \text{space } M. \text{norm } (s \ (r \ i) \ y) \neq 0\} \neq \infty \ \bigwedge x. x \in \text{space } M \Rightarrow (\lambda i. \text{norm } (s \ (r \ i) \ x)) \longrightarrow \text{norm } (f \ x) \ \bigwedge i. x \in \text{space } M \Rightarrow \text{norm } (\text{norm } (s \ (r \ i) \ x)) \leq 2 * \text{norm } (\text{norm } (f \ x))$

using *s* **by** (*auto intro: LIMSEQ-subseq-LIMSEQ*[*OF tendsto-norm r(2)*, *unfolded comp-def*] *simple-function-compose1*)

obtain *r'* **where** *r'*: *AE x in M. (λi. (cond-exp M F (λx. norm (s (r' i)) x)) x) ⟶ cond-exp M F (λx. norm (f x)) x strict-mono r'* **using** *cond-exp-simple-lim*[*OF integrable-norm norm-s-r, OF assms*] **by** *blast*

have *AE x in M. ∀ i. norm (cond-exp M F (s (r' i))) x ≤ cond-exp M F (λx. norm (s (r' i)) x) x* **using** *s* **by** (*auto intro: cond-exp-contraction-simple simp add: AE-all-countable*)

moreover **have** *AE x in M. (λi. norm (cond-exp M F (s (r' i))) x) ⟶ norm (cond-exp M F f x)* **using** *r* *LIMSEQ-subseq-LIMSEQ*[*OF tendsto-norm r'(2)*, *unfolded comp-def*] **by** *fast*

ultimately show *?thesis* **using** *LIMSEQ-le r'(1)* **by** *fast*
qed

lemma *cond-exp-measurable-mult*:

fixes *f g :: 'a ⇒ real*

assumes [*measurable*]: *integrable M (λx. f x * g x) integrable M g f ∈ borel-measurable F*

shows *integrable M (λx. f x * cond-exp M F g x)*

*AE x in M. cond-exp M F (λx. f x * g x) x = f x * cond-exp M F g x*

proof–

show *integrable: integrable M (λx. f x * cond-exp M F g x)* **using** *cond-exp-real*[*OF assms(2)*] **by** (*intro integrable-cong-AE-imp*[*OF real-cond-exp-intg(1)*, *OF assms(1,3)*] *assms(2)*[*THEN borel-measurable-integrable*]) *measurable-from-subalg*[*OF subalg*])

auto

interpret *sigma-finite-measure restr-to-subalg M F* **by** (*rule sigma-fin-subalg*)

{

fix *A* **assume** *asm: A ∈ sets F*

hence *asm'*: *A ∈ sets M* **using** *subalg* **by** (*fastforce simp add: subalgebra-def*)

have *set-lebesgue-integral M A (cond-exp M F (λx. f x * g x)) = set-lebesgue-integral M A (λx. f x * g x)* **by** (*simp add: cond-exp-set-integral*[*OF assms(1) asm*])

also **have** ... = *set-lebesgue-integral M A (λx. f x * real-cond-exp M F g x)* **using** *borel-measurable-times*[*OF borel-measurable-indicator*[*OF asm*] *assms(3)*] *borel-measurable-integrable*[*OF assms(2)*] *integrable-mult-indicator*[*OF asm' assms(1)*]

by (*fastforce simp add: set-lebesgue-integral-def mult.assoc[symmetric] intro: real-cond-exp-intg(2)[symmetric]*)

also have ... = *set-lebesgue-integral* $M A (\lambda x. f x * \text{cond-exp } M F g x)$ **using** *cond-exp-real*[$OF \text{ assms}(2)$] *asm'* *borel-measurable-cond-exp'* *borel-measurable-cond-exp2* *measurable-from-subalg*[$OF \text{ subalg assms}(3)$] **by** (*auto simp add: set-lebesgue-integral-def intro: integral-cong-AE*)

finally have *set-lebesgue-integral* $M A (\text{cond-exp } M F (\lambda x. f x * g x)) = \int_{x \in A}. (f x * \text{cond-exp } M F g x) \partial M$.
 $\}$

hence $AE x \text{ in } \text{restr-to-subalg } M F. \text{cond-exp } M F (\lambda x. f x * g x) x = f x * \text{cond-exp } M F g x$ **by** (*intro density-unique integrable-cond-exp integrable integrable-in-subalg subalg, measurable, simp add: set-lebesgue-integral-def integral-subalgebra2*[$OF \text{ subalg}$] *sets-restr-to-subalg*[$OF \text{ subalg}$])

thus $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x * g x) x = f x * \text{cond-exp } M F g x$ **by** (*rule AE-restr-to-subalg*[$OF \text{ subalg}$])

qed

lemma *cond-exp-measurable-scaleR*:

fixes $f :: 'a \Rightarrow \text{real}$ **and** $g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$

assumes [*measurable*]: *integrable* $M (\lambda x. f x *_R g x)$ *integrable* $M g f \in \text{borel-measurable } F$

shows *integrable* $M (\lambda x. f x *_R \text{cond-exp } M F g x)$

$AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x *_R g x) x = f x *_R \text{cond-exp } M F g x$

proof –

let $?F = \text{restr-to-subalg } M F$

have *subalg'*: *subalgebra* $M (\text{restr-to-subalg } M F)$ **by** (*metis sets-eq-imp-space-eq sets-restr-to-subalg subalg subalgebra-def*)

$\{$

fix z **assume** *asm*[*measurable*]: *integrable* $M (\lambda x. z x *_R g x)$ $z \in \text{borel-measurable } ?F$

hence *asm'*[*measurable*]: $z \in \text{borel-measurable } F$ **using** *measurable-in-subalg'* *subalg* **by** *blast*

have *integrable* $M (\lambda x. z x *_R \text{cond-exp } M F g x)$ *LINT* $x | M. z x *_R g x = \text{LINT } x | M. z x *_R \text{cond-exp } M F g x$

proof –

obtain s **where** $s\text{-is}$: $\bigwedge i. \text{simple-function } ?F (s i) \bigwedge x. x \in \text{space } ?F \implies (\lambda i. s i x) \longrightarrow z x \bigwedge i x. x \in \text{space } ?F \implies \text{norm } (s i x) \leq 2 * \text{norm } (z x)$ **using** *borel-measurable-implies-sequence-metric*[$OF \text{ asm}(2), \text{of } 0$] **by** *force*

have $s\text{-scaleR-g-tendsto}$: $AE x \text{ in } M. (\lambda i. s i x *_R g x) \longrightarrow z x *_R g x$ **using** $s\text{-is}(2)$ **by** (*simp add: space-restr-to-subalg tendsto-scaleR*)

have $s\text{-scaleR-cond-exp-g-tendsto}$: $AE x \text{ in } ?F. (\lambda i. s i x *_R \text{cond-exp } M F g x) \longrightarrow z x *_R \text{cond-exp } M F g x$ **using** $s\text{-is}(2)$ **by** (*simp add: tendsto-scaleR*)

have $s\text{-scaleR-g-meas}$: $(\lambda x. s i x *_R g x) \in \text{borel-measurable } M$ **for** i **using** $s\text{-is}(1)$ [*THEN borel-measurable-simple-function, THEN subalg'*[*THEN measurable-from-subalg*]] **by** *simp*

have $s\text{-scaleR-cond-exp-g-meas}$: $(\lambda x. s i x *_R \text{cond-exp } M F g x) \in \text{borel-measurable } ?F$ **for** i **using** $s\text{-is}(1)$ [*THEN borel-measurable-simple-function*] *measurable-in-subalg*[OF

subalg borel-measurable-cond-exp **by** (*fastforce intro: borel-measurable-scaleR*)

have *s-scaleR-g-AE-bdd*: *AE x in M. norm (s i x *_R g x) ≤ 2 * norm (z x *_R g x)* **for** *i* **using** *s-is(3)* **by** (*fastforce simp add: space-restr-to-subalg mult.assoc[symmetric] mult-right-mono*)

{

fix *i*

have *asm: integrable M (λx. norm (z x) * norm (g x))* **using** *asm(1)* [*THEN integrable-norm*] **by** *simp*

have *AE x in ?F. norm (s i x *_R cond-exp M F g x) ≤ 2 * norm (z x) * norm (cond-exp M F g x)* **using** *s-is(3)* **by** (*fastforce simp add: mult-mono*)

moreover **have** *AE x in ?F. norm (z x) * cond-exp M F (λx. norm (g x)) x = cond-exp M F (λx. norm (z x) * norm (g x)) x* **by** (*rule cond-exp-measurable-mult(2)*) [*THEN AE-symmetric, OF asm integrable-norm, OF assms(2), THEN AE-restr-to-subalg2*] [*OF subalg*], *auto*)

ultimately **have** *AE x in ?F. norm (s i x *_R cond-exp M F g x) ≤ 2 * cond-exp M F (λx. norm (z x *_R g x)) x* **using** *cond-exp-contraction* [*OF assms(2)*], *THEN AE-restr-to-subalg2* [*OF subalg*] *order-trans* [*OF - mult-mono*] **by** *fastforce*

}

note *s-scaleR-cond-exp-g-AE-bdd = this*

{

fix *i*

have *s-meas-M[measurable]: s i ∈ borel-measurable M* **by** (*meson borel-measurable-simple-function measurable-from-subalg s-is(1) subalg'*)

have *s-meas-F[measurable]: s i ∈ borel-measurable F* **by** (*meson borel-measurable-simple-function measurable-in-subalg' s-is(1) subalg*)

have *s-scaleR-eq: s i x *_R h x = (∑ y ∈ s i ' space M. (indicator (s i - ' {y} ∩ space M) x *_R y) *_R h x)* **if** *x ∈ space M* **for** *x* **and** *h :: 'a ⇒ 'b* **using** *simple-function-indicator-representation* [*OF s-is(1), of x i*] **that** *unfolding space-restr-to-subalg scaleR-left.sum* [*of - - h x, symmetric*] **by** *presburger*

have *LINT x|M. s i x *_R g x = LINT x|M. (∑ y ∈ s i ' space M. indicator (s i - ' {y} ∩ space M) x *_R y *_R g x)* **using** *s-scaleR-eq* **by** (*intro Bochner-Integration.integral-cong*) *auto*

also **have** *... = (∑ y ∈ s i ' space M. LINT x|M. indicator (s i - ' {y} ∩ space M) x *_R y *_R g x)* **by** (*intro Bochner-Integration.integral-sum integrable-mult-indicator* [*OF - integrable-scaleR-right*] *assms(2)*) *simp*

also **have** *... = (∑ y ∈ s i ' space M. y *_R set-lebesgue-integral M (s i - ' {y} ∩ space M) g)* **by** (*simp only: set-lebesgue-integral-def[symmetric]*) *simp*

also **have** *... = (∑ y ∈ s i ' space M. y *_R set-lebesgue-integral M (s i - ' {y} ∩ space M) (cond-exp M F g))* **using** *assms(2)* *subalg borel-measurable-vimage* [*OF s-meas-F*] **by** (*subst cond-exp-set-integral, auto simp add: subalgebra-def*)

also **have** *... = (∑ y ∈ s i ' space M. LINT x|M. indicator (s i - ' {y} ∩ space M) x *_R y *_R cond-exp M F g x)* **by** (*simp only: set-lebesgue-integral-def[symmetric]*) *simp*

also have ... = $LINT\ x|M. (\sum_{y \in s\ i\ 'space\ M.} indicator\ (s\ i\ -'\{y\} \cap space\ M)\ x *_R y *_R cond-exp\ M\ F\ g\ x)$ **by** (intro Bochner-Integration.integral-sum[symmetric] integrable-mult-indicator[OF - integrable-scaleR-right]) auto
also have ... = $LINT\ x|M. s\ i\ x *_R cond-exp\ M\ F\ g\ x$ **using** s-scaleR-eq
by (intro Bochner-Integration.integral-cong) auto
finally have $LINT\ x|M. s\ i\ x *_R g\ x = LINT\ x|?F. s\ i\ x *_R cond-exp\ M\ F\ g\ x$ **by** (simp add: integral-subalgebra2[OF subalg])
}
note integral-s-eq = this

show integrable $M\ (\lambda x. z\ x *_R cond-exp\ M\ F\ g\ x)$ **using** s-scaleR-cond-exp-g-meas asm(2) borel-measurable-cond-exp' **by** (intro integrable-from-subalg[OF subalg] integrable-cond-exp integrable-dominated-convergence[OF - - - s-scaleR-cond-exp-g-tendsto s-scaleR-cond-exp-g-AE-bdd]) (auto intro: measurable-from-subalg[OF subalg] integrable-in-subalg measurable-in-subalg subalg)

have $(\lambda i. LINT\ x|M. s\ i\ x *_R g\ x) \longrightarrow LINT\ x|M. z\ x *_R g\ x$ **using** s-scaleR-g-meas asm(1)[THEN integrable-norm] asm' borel-measurable-cond-exp' **by** (intro integral-dominated-convergence[OF - - - s-scaleR-g-tendsto s-scaleR-g-AE-bdd]) (auto intro: measurable-from-subalg[OF subalg])

moreover have $(\lambda i. LINT\ x|?F. s\ i\ x *_R cond-exp\ M\ F\ g\ x) \longrightarrow LINT\ x|?F. z\ x *_R cond-exp\ M\ F\ g\ x$ **using** s-scaleR-cond-exp-g-meas asm(2) borel-measurable-cond-exp' **by** (intro integral-dominated-convergence[OF - - - s-scaleR-cond-exp-g-tendsto s-scaleR-cond-exp-g-AE-bdd]) (auto intro: measurable-from-subalg[OF subalg] integrable-in-subalg measurable-in-subalg subalg)

ultimately show $LINT\ x|M. z\ x *_R g\ x = LINT\ x|M. z\ x *_R cond-exp\ M\ F\ g\ x$ **using** integral-s-eq **using** subalg **by** (simp add: LIMSEQ-unique integral-subalgebra2)

qed

}

note * = this

show integrable $M\ (\lambda x. f\ x *_R cond-exp\ M\ F\ g\ x)$ **using** * assms measurable-in-subalg[OF subalg] **by** blast

{

fix A **assume** $asm: A \in F$

hence integrable $M\ (\lambda x. indicat-real\ A\ x *_R f\ x *_R g\ x)$ **using** subalg **by** (fastforce simp add: subalgebra-def intro!: integrable-mult-indicator assms(1))

hence set-lebesgue-integral $M\ A\ (\lambda x. f\ x *_R g\ x) = set-lebesgue-integral\ M\ A\ (\lambda x. f\ x *_R cond-exp\ M\ F\ g\ x)$ **unfolding** set-lebesgue-integral-def **using** asm **by** (auto intro!: * measurable-in-subalg[OF subalg])

}

thus $AE\ x\ in\ M. cond-exp\ M\ F\ (\lambda x. f\ x *_R g\ x)\ x = f\ x *_R cond-exp\ M\ F\ g\ x$ **using** borel-measurable-cond-exp **by** (intro cond-exp-charact, auto intro!: * assms)

measurable-in-subalg[*OF subalg*])
qed

lemma *cond-exp-sum* [*intro, simp*]:
fixes $f :: 't \Rightarrow 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes [*measurable*]: $\bigwedge i. \text{integrable } M (f i)$
shows $AE\ x\ \text{in } M. \text{cond-exp } M\ F\ (\lambda x. \sum_{i \in I}. f\ i\ x)\ x = (\sum_{i \in I}. \text{cond-exp } M\ F\ (f\ i)\ x)$
proof (*rule has-cond-exp-charact, intro has-cond-expI'*)
fix A **assume** [*measurable*]: $A \in \text{sets } F$
then have $A\text{-meas } [measurable]: A \in \text{sets } M$ **by** (*meson subsetD subalg subalgebra-def*)

have $(\int x \in A. (\sum_{i \in I}. f\ i\ x) \partial M) = (\int x. (\sum_{i \in I}. \text{indicator } A\ x\ *_R\ f\ i\ x) \partial M)$
unfolding *set-lebesgue-integral-def* **by** (*simp add: scaleR-sum-right*)
also have $\dots = (\sum_{i \in I}. (\int x. \text{indicator } A\ x\ *_R\ f\ i\ x\ \partial M))$ **using** *assms* **by** (*auto intro!: Bochner-Integration.integral-sum integrable-mult-indicator*)
also have $\dots = (\sum_{i \in I}. (\int x. \text{indicator } A\ x\ *_R\ \text{cond-exp } M\ F\ (f\ i)\ x\ \partial M))$ **using** *cond-exp-set-integral*[*OF assms*] **by** (*simp add: set-lebesgue-integral-def*)
also have $\dots = (\int x. (\sum_{i \in I}. \text{indicator } A\ x\ *_R\ \text{cond-exp } M\ F\ (f\ i)\ x) \partial M)$
using *assms* **by** (*auto intro!: Bochner-Integration.integral-sum[symmetric] integrable-mult-indicator*)
also have $\dots = (\int x \in A. (\sum_{i \in I}. \text{cond-exp } M\ F\ (f\ i)\ x) \partial M)$ **unfolding** *set-lebesgue-integral-def*
by (*simp add: scaleR-sum-right*)
finally show $(\int x \in A. (\sum_{i \in I}. f\ i\ x) \partial M) = (\int x \in A. (\sum_{i \in I}. \text{cond-exp } M\ F\ (f\ i)\ x) \partial M)$ **by** *auto*
qed (*auto simp add: assms integrable-cond-exp*)

6.1 Linearly Ordered Banach Spaces

lemma *cond-exp-gr-c*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}, \text{linorder-topology}, \text{ordered-real-vector}\}$
assumes *integrable* $M\ f\ AE\ x\ \text{in } M. f\ x > c$
shows $AE\ x\ \text{in } M. \text{cond-exp } M\ F\ f\ x > c$
proof –
define X **where** $X = \{x \in \text{space } M. \text{cond-exp } M\ F\ f\ x \leq c\}$
have [*measurable*]: $X \in \text{sets } F$ **unfolding** $X\text{-def}$ **by** *measurable* (*metis sets.top subalg subalgebra-def*)
hence $X\text{-in-}M: X \in \text{sets } M$ **using** *sets-restr-to-subalg subalg subalgebra-def* **by** *blast*
have $\text{emeasure } M\ X = 0$
proof (*rule ccontr*)
assume $\text{emeasure } M\ X \neq 0$
have $\text{emeasure } (restr\text{-to-subalg } M\ F)\ X = \text{emeasure } M\ X$ **by** (*simp add: emeasure-restr-to-subalg subalg*)
hence $\text{emeasure } (restr\text{-to-subalg } M\ F)\ X > 0$ **using** $\neg(\text{emeasure } M\ X) = 0$
gr-zeroI **by** *auto*
then obtain A **where** $A: A \in \text{sets } (restr\text{-to-subalg } M\ F)\ A \subseteq X$ emeasure

$(\text{restr-to-subalg } M \ F) \ A > 0 \ \text{emeasure } (\text{restr-to-subalg } M \ F) \ A < \infty$
using *sigma-fin-subalg* **by** (*metis emeasure-notin-sets ennreal-0 infinity-ennreal-def*
le-less-linear neq-top-trans not-gr-zero order-refl sigma-finite-measure.approx-PInf-emeasure-with-finite)
hence $[simp]: A \in \text{sets } F$ **using** *subalg sets-restr-to-subalg* **by** *blast*
hence $A\text{-in-sets-}M[simp]: A \in \text{sets } M$ **using** *sets-restr-to-subalg subalg subalgebra-def* **by** *blast*
have $[simp]: \text{set-integrable } M \ A \ (\lambda x. c)$ **using** *A subalg* **by** (*auto simp add: set-integrable-def emeasure-restr-to-subalg*)
have $[simp]: \text{set-integrable } M \ A \ f$ **unfolding** *set-integrable-def* **by** (*rule integrable-mult-indicator, auto simp add: assms(1)*)
have $AE \ x \text{ in } M. \text{indicator } A \ x *_R c = \text{indicator } A \ x *_R f \ x$
proof (*rule integral-eq-mono-AE-eq-AE*)
show $LINT \ x|M. \text{indicator } A \ x *_R c = LINT \ x|M. \text{indicator } A \ x *_R f \ x$
proof (*simp only: set-lebesgue-integral-def[symmetric], rule antisym*)
show $(\int x \in A. c \ \partial M) \leq (\int x \in A. f \ x \ \partial M)$ **using** *assms(2)* **by** (*intro set-integral-mono-AE-banach*) *auto*
have $(\int x \in A. f \ x \ \partial M) = (\int x \in A. \text{cond-exp } M \ F \ f \ x \ \partial M)$ **by** (*rule cond-exp-set-integral, auto simp add: assms*)
also have $\dots \leq (\int x \in A. c \ \partial M)$ **using** *A* **by** (*auto intro!: set-integral-mono-banach simp add: X-def*)
finally show $(\int x \in A. f \ x \ \partial M) \leq (\int x \in A. c \ \partial M)$ **by** *simp*
qed
show $AE \ x \text{ in } M. \text{indicator } A \ x *_R c \leq \text{indicator } A \ x *_R f \ x$ **using** *assms* **by** (*auto simp add: X-def indicator-def*)
qed (*auto simp add: set-integrable-def[symmetric]*)
hence $AE \ x \in A \text{ in } M. c = f \ x$ **by** *auto*
hence $AE \ x \in A \text{ in } M. \text{False}$ **using** *assms(2)* **by** *auto*
hence $A \in \text{null-sets } M$ **using** *AE-iff-null-sets A-in-sets-M* **by** *metis*
thus False **using** *A(3)* **by** (*simp add: emeasure-restr-to-subalg null-setsD1 subalg*)
qed
thus *?thesis* **using** *AE-iff-null-sets[OF X-in-M]* **unfolding** *X-def* **by** *auto*
qed

corollary *cond-exp-less-c*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes $\text{integrable } M \ f \ AE \ x \text{ in } M. f \ x < c$

shows $AE \ x \text{ in } M. \text{cond-exp } M \ F \ f \ x < c$

proof –

have $AE \ x \text{ in } M. \text{cond-exp } M \ F \ f \ x = - \text{cond-exp } M \ F \ (\lambda x. - f \ x) \ x$ **using** *cond-exp-uminus[OF assms(1)]* **by** *auto*

moreover have $AE \ x \text{ in } M. \text{cond-exp } M \ F \ (\lambda x. - f \ x) \ x > - c$ **using** *assms* **by** (*intro cond-exp-gr-c*) *auto*

ultimately show *?thesis* **by** (*force simp add: minus-less-iff*)

qed

lemma *cond-exp-mono-strict*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$

```

dered-real-vector}
  assumes integrable M f integrable M g AE x in M. f x < g x
  shows AE x in M. cond-exp M F f x < cond-exp M F g x
  using cond-exp-less-c[OF Bochner-Integration.integrable-diff, OF assms(1,2), of
0]
    cond-exp-diff[OF assms(1,2)] assms(3) by auto

lemma cond-exp-ge-c:
  fixes f :: 'a  $\Rightarrow$  'b :: {second-countable-topology, banach, linorder-topology, or-
dered-real-vector}
  assumes [measurable]: integrable M f
    and AE x in M. f x  $\geq$  c
  shows AE x in M. cond-exp M F f x  $\geq$  c
proof -
  let ?F = restr-to-subalg M F
  interpret sigma-finite-measure restr-to-subalg M F using sigma-fin-subalg by
auto
  {
    fix A assume asm: A  $\in$  sets ?F 0 < measure ?F A
    have [simp]: sets ?F = sets F measure ?F A = measure M A using asm by (auto
simp add: measure-def sets-restr-to-subalg[OF subalg] emeasure-restr-to-subalg[OF
subalg])
    have M-A: emeasure M A <  $\infty$  using measure-zero-top asm by (force simp
add: top.not-eq-extremum)
    hence F-A: emeasure ?F A <  $\infty$  using asm(1) emeasure-restr-to-subalg subalg
by fastforce
    have set-lebesgue-integral M A ( $\lambda$ -. c)  $\leq$  set-lebesgue-integral M A f using
assms asm M-A subalg by (intro set-integral-mono-AE-banach, auto simp add:
set-integrable-def integrable-mult-indicator subalgebra-def sets-restr-to-subalg)
    also have ... = set-lebesgue-integral M A (cond-exp M F f) using cond-exp-set-integral[OF
assms(1)] asm by auto
    also have ... = set-lebesgue-integral ?F A (cond-exp M F f) unfolding set-lebesgue-integral-def
using asm borel-measurable-cond-exp by (intro integral-subalgebra2[OF subalg, sym-
metric], simp)
    finally have (1 / measure ?F A) *R set-lebesgue-integral ?F A (cond-exp M F f)
 $\in$  {c..} using asm subalg M-A by (auto simp add: set-integral-const subalgebra-def
intro!: pos-divideR-le-eq[THEN iffD1])
  }
  thus ?thesis using AE-restr-to-subalg[OF subalg] averaging-theorem[OF inte-
grable-in-subalg closed-atLeast, OF subalg borel-measurable-cond-exp integrable-cond-exp]
by auto
qed

corollary cond-exp-le-c:
  fixes f :: 'a  $\Rightarrow$  'b :: {second-countable-topology, banach, linorder-topology, or-
dered-real-vector}
  assumes integrable M f
    and AE x in M. f x  $\leq$  c
  shows AE x in M. cond-exp M F f x  $\leq$  c

```

proof –

have $AE\ x\ in\ M. \text{cond-exp}\ M\ F\ f\ x = -\ \text{cond-exp}\ M\ F\ (\lambda x. -\ f\ x)\ x$ **using** $\text{cond-exp-uminus}[OF\ \text{assms}(1)]$ **by** *force*
moreover **have** $AE\ x\ in\ M. \text{cond-exp}\ M\ F\ (\lambda x. -\ f\ x)\ x \geq -\ c$ **using** *assms*
by $(\text{intro}\ \text{cond-exp-ge-c})\ \text{auto}$
ultimately show *?thesis* **by** $(\text{force}\ \text{simp}\ \text{add:}\ \text{minus-le-iff})$
qed

corollary *cond-exp-mono*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes $\text{integrable}\ M\ f\ \text{integrable}\ M\ g\ AE\ x\ in\ M. f\ x \leq g\ x$
shows $AE\ x\ in\ M. \text{cond-exp}\ M\ F\ f\ x \leq \text{cond-exp}\ M\ F\ g\ x$
using $\text{cond-exp-le-c}[OF\ \text{Bochner-Integration.integrable-diff}, OF\ \text{assms}(1,2),\ of\ 0]$
 $\text{cond-exp-diff}[OF\ \text{assms}(1,2)]\ \text{assms}(3)$ **by** *auto*

corollary *cond-exp-min*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes $\text{integrable}\ M\ f\ \text{integrable}\ M\ g$
shows $AE\ \xi\ in\ M. \text{cond-exp}\ M\ F\ (\lambda x. \min (f\ x)\ (g\ x))\ \xi \leq \min (\text{cond-exp}\ M\ F\ f\ \xi)\ (\text{cond-exp}\ M\ F\ g\ \xi)$
proof –
have $AE\ \xi\ in\ M. \text{cond-exp}\ M\ F\ (\lambda x. \min (f\ x)\ (g\ x))\ \xi \leq \text{cond-exp}\ M\ F\ f\ \xi$ **by** $(\text{intro}\ \text{cond-exp-mono}\ \text{integrable-min}\ \text{assms},\ \text{simp})$
moreover **have** $AE\ \xi\ in\ M. \text{cond-exp}\ M\ F\ (\lambda x. \min (f\ x)\ (g\ x))\ \xi \leq \text{cond-exp}\ M\ F\ g\ \xi$ **by** $(\text{intro}\ \text{cond-exp-mono}\ \text{integrable-min}\ \text{assms},\ \text{simp})$
ultimately show $AE\ \xi\ in\ M. \text{cond-exp}\ M\ F\ (\lambda x. \min (f\ x)\ (g\ x))\ \xi \leq \min (\text{cond-exp}\ M\ F\ f\ \xi)\ (\text{cond-exp}\ M\ F\ g\ \xi)$ **by** *fastforce*
qed

corollary *cond-exp-max*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes $\text{integrable}\ M\ f\ \text{integrable}\ M\ g$
shows $AE\ \xi\ in\ M. \text{cond-exp}\ M\ F\ (\lambda x. \max (f\ x)\ (g\ x))\ \xi \geq \max (\text{cond-exp}\ M\ F\ f\ \xi)\ (\text{cond-exp}\ M\ F\ g\ \xi)$
proof –
have $AE\ \xi\ in\ M. \text{cond-exp}\ M\ F\ (\lambda x. \max (f\ x)\ (g\ x))\ \xi \geq \text{cond-exp}\ M\ F\ f\ \xi$ **by** $(\text{intro}\ \text{cond-exp-mono}\ \text{integrable-max}\ \text{assms},\ \text{simp})$
moreover **have** $AE\ \xi\ in\ M. \text{cond-exp}\ M\ F\ (\lambda x. \max (f\ x)\ (g\ x))\ \xi \geq \text{cond-exp}\ M\ F\ g\ \xi$ **by** $(\text{intro}\ \text{cond-exp-mono}\ \text{integrable-max}\ \text{assms},\ \text{simp})$
ultimately show $AE\ \xi\ in\ M. \text{cond-exp}\ M\ F\ (\lambda x. \max (f\ x)\ (g\ x))\ \xi \geq \max (\text{cond-exp}\ M\ F\ f\ \xi)\ (\text{cond-exp}\ M\ F\ g\ \xi)$ **by** *fastforce*
qed

corollary *cond-exp-inf*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$

```

dered-real-vector, lattice}
  assumes integrable M f integrable M g
  shows  $\text{AE } \xi \text{ in } M. \text{ cond-exp } M F (\lambda x. \inf (f x) (g x)) \xi \leq \inf (\text{cond-exp } M F f \xi) (\text{cond-exp } M F g \xi)$ 
  unfolding inf-min using assms by (rule cond-exp-min)

corollary cond-exp-sup:
  fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector, lattice}\}$ 
  assumes integrable M f integrable M g
  shows  $\text{AE } \xi \text{ in } M. \text{ cond-exp } M F (\lambda x. \sup (f x) (g x)) \xi \geq \sup (\text{cond-exp } M F f \xi) (\text{cond-exp } M F g \xi)$ 
  unfolding sup-max using assms by (rule cond-exp-max)

end

end

theory Filtered-Measure
imports HOL-Probability.Conditional-Expectation
begin

```

7 Filtered Measure Spaces

7.1 Filtered Measure

```

locale filtered-measure =
  fixes M F and  $t_0 :: 'b :: \{\text{second-countable-topology, linorder-topology}\}$ 
  assumes subalgebra:  $\bigwedge i. t_0 \leq i \implies \text{subalgebra } M (F i)$ 
  and sets-F-mono:  $\bigwedge i j. t_0 \leq i \implies i \leq j \implies \text{sets } (F i) \subseteq \text{sets } (F j)$ 
begin

lemma space-F:
  assumes  $t_0 \leq i$ 
  shows  $\text{space } (F i) = \text{space } M$ 
  using subalgebra assms by (simp add: subalgebra-def)

lemma subalgebra-F:
  assumes  $t_0 \leq i \leq j$ 
  shows  $\text{subalgebra } (F j) (F i)$ 
  unfolding subalgebra-def using assms by (simp add: space-F sets-F-mono)

lemma borel-measurable-mono:
  assumes  $t_0 \leq i \leq j$ 
  shows  $\text{borel-measurable } (F i) \subseteq \text{borel-measurable } (F j)$ 
  unfolding subset-iff by (metis assms subalgebra-F measurable-from-subalg)

end

```

locale *nat-filtered-measure* = *filtered-measure* *M F 0* **for** *M* **and** *F :: nat* \Rightarrow -
locale *real-filtered-measure* = *filtered-measure* *M F 0* **for** *M* **and** *F :: real* \Rightarrow -

context *nat-filtered-measure*
begin

lemma *space-F*: *space (F i) = space M*
using *subalgebra* **by** (*simp add: subalgebra-def*)

lemma *subalgebra-F*:
assumes *i* \leq *j*
shows *subalgebra (F j) (F i)*
unfolding *subalgebra-def* **using** *assms* **by** (*simp add: space-F sets-F-mono*)

lemma *borel-measurable-mono*:
assumes *i* \leq *j*
shows *borel-measurable (F i) \subseteq borel-measurable (F j)*
unfolding *subset-iff* **by** (*metis assms subalgebra-F measurable-from-subalg*)

end

7.2 Sigma Finite Filtered Measure

The locale presented here is a generalization of the *sigma-finite-subalgebra* for a particular filtration.

locale *sigma-finite-filtered-measure* = *filtered-measure* +
assumes *sigma-finite*: *sigma-finite-subalgebra M (F t₀)*

lemma (**in** *sigma-finite-filtered-measure*) *sigma-finite-subalgebra-F[intro]*:
assumes *t₀* \leq *i*
shows *sigma-finite-subalgebra M (F i)*
using *assms* **by** (*metis dual-order.refl sets-F-mono sigma-finite sigma-finite-subalgebra.nested-subalg-is-sigma-subalgebra subalgebra-def*)

locale *nat-sigma-finite-filtered-measure* = *sigma-finite-filtered-measure M F 0* **for** *M* **and** *F :: nat* \Rightarrow -

locale *real-sigma-finite-filtered-measure* = *sigma-finite-filtered-measure M F 0* **for** *M* **and** *F :: real* \Rightarrow -

sublocale *nat-sigma-finite-filtered-measure* \subseteq *nat-filtered-measure* ..
sublocale *nat-sigma-finite-filtered-measure* \subseteq *sigma-finite-subalgebra M F i* **by** *blast*

sublocale *real-sigma-finite-filtered-measure* \subseteq *real-filtered-measure* ..

7.3 Filtered Finite Measure

locale *finite-filtered-measure* = *filtered-measure* + *finite-measure*

sublocale *finite-filtered-measure* \subseteq *sigma-finite-filtered-measure*
using *subalgebra* **by** (*unfold-locales*, *blast*, *meson* *dual-order.refl* *finite-measure-axioms*
finite-measure-def *finite-measure-restr-to-subalg* *sigma-finite-measure.sigma-finite-countable*
subalgebra)

locale *nat-finite-filtered-measure* = *finite-filtered-measure* *M F 0* **for** *M* **and** *F* ::
nat \Rightarrow -
locale *real-finite-filtered-measure* = *finite-filtered-measure* *M F 0* **for** *M* **and** *F* ::
real \Rightarrow -

sublocale *nat-finite-filtered-measure* \subseteq *nat-sigma-finite-filtered-measure* ..
sublocale *real-finite-filtered-measure* \subseteq *real-sigma-finite-filtered-measure* ..

7.4 Constant Filtration

lemma *filtered-measure-constant-filtration*:

assumes *subalgebra* *M F*
shows *filtered-measure* *M* (λ -. *F*) *t*₀
using *assms* **by** (*unfold-locales*) (*auto simp add: subalgebra-def*)

sublocale *sigma-finite-subalgebra* \subseteq *constant-filtration: sigma-finite-filtered-measure*
M λ -. :: '*t* :: {*second-countable-topology*, *linorder-topology*}. *F* *t*₀
using *subalg* **by** (*unfold-locales*) (*auto simp add: subalgebra-def*)

end
theory *Stochastic-Process*
imports *Filtered-Measure* *Measure-Space-Addendum*
begin

8 Stochastic Processes

8.1 Stochastic Process

A stochastic process is a collection of random variables, indexed by a type '*b*.

locale *stochastic-process* =
fixes *M* *t*₀ **and** *X* :: '*b* :: {*second-countable-topology*, *linorder-topology*} \Rightarrow '*a* \Rightarrow
'*c* :: {*second-countable-topology*, *banach*}
assumes *random-variable[measurable]*: $\bigwedge i. t_0 \leq i \implies X\ i \in \text{borel-measurable } M$
begin

definition *left-continuous* **where** *left-continuous* = (*AE* ξ *in* *M*. $\forall t. \text{continuous}$
(*at-left* *t*) ($\lambda i. X\ i\ \xi$))

definition *right-continuous* **where** *right-continuous* = (*AE* ξ *in* *M*. $\forall t. \text{continuous}$
(*at-right* *t*) ($\lambda i. X\ i\ \xi$))

end

locale *nat-stochastic-process* = *stochastic-process* $M\ 0 :: \text{nat } X$ **for** $M\ X$
locale *real-stochastic-process* = *stochastic-process* $M\ 0 :: \text{real } X$ **for** $M\ X$

lemma *stochastic-process-const-fun*:
assumes $f \in \text{borel-measurable } M$
shows *stochastic-process* $M\ t_0\ (\lambda\ -. \ f)$ **using** *assms* **by** (*unfold-locale*)

lemma *stochastic-process-const*:
shows *stochastic-process* $M\ t_0\ (\lambda\ i\ -. \ c\ i)$ **by** (*unfold-locale*) *simp*

context *stochastic-process*
begin

lemma *compose*:
assumes $\bigwedge i. t_0 \leq i \implies f\ i \in \text{borel-measurable } \text{borel}$
shows *stochastic-process* $M\ t_0\ (\lambda\ i\ \xi. (f\ i)\ (X\ i\ \xi))$
by (*unfold-locale*) (*intro measurable-compose[OF random-variable assms]*)

lemma *norm*: *stochastic-process* $M\ t_0\ (\lambda\ i\ \xi. \text{norm } (X\ i\ \xi))$ **by** (*fastforce intro: compose*)

lemma *scaleR-right*:
assumes *stochastic-process* $M\ t_0\ Y$
shows *stochastic-process* $M\ t_0\ (\lambda\ i\ \xi. (Y\ i\ \xi) *_{\mathbb{R}} (X\ i\ \xi))$
using *stochastic-process.random-variable[OF assms]* *random-variable* **by** (*unfold-locale*) *simp*

lemma *scaleR-right-const-fun*:
assumes $f \in \text{borel-measurable } M$
shows *stochastic-process* $M\ t_0\ (\lambda\ i\ \xi. f\ \xi *_{\mathbb{R}} (X\ i\ \xi))$
by (*unfold-locale*) (*intro borel-measurable-scaleR assms random-variable*)

lemma *scaleR-right-const*: *stochastic-process* $M\ t_0\ (\lambda\ i\ \xi. c\ i *_{\mathbb{R}} (X\ i\ \xi))$
by (*unfold-locale*) *simp*

lemma *add*:
assumes *stochastic-process* $M\ t_0\ Y$
shows *stochastic-process* $M\ t_0\ (\lambda\ i\ \xi. X\ i\ \xi + Y\ i\ \xi)$
using *stochastic-process.random-variable[OF assms]* *random-variable* **by** (*unfold-locale*) *simp*

lemma *diff*:
assumes *stochastic-process* $M\ t_0\ Y$
shows *stochastic-process* $M\ t_0\ (\lambda\ i\ \xi. X\ i\ \xi - Y\ i\ \xi)$
using *stochastic-process.random-variable[OF assms]* *random-variable* **by** (*unfold-locale*) *simp*

lemma *uminus*: *stochastic-process* $M\ t_0\ (-X)$ **using** *scaleR-right-const[of $\lambda\ -. \ -1$]*
by (*simp add: fun-Compl-def*)

lemma *partial-sum*: *stochastic-process* $M\ t_0\ (\lambda n\ \xi.\ \sum_{i \in \{t_0..<n\}}. X\ i\ \xi)$ **by** (*unfold-locales*) *simp*

lemma *partial-sum'*: *stochastic-process* $M\ t_0\ (\lambda n\ \xi.\ \sum_{i \in \{t_0..n\}}. X\ i\ \xi)$ **by** (*unfold-locales*) *simp*

end

lemma *stochastic-process-sum*:

assumes $\bigwedge i. i \in I \implies \text{stochastic-process } M\ t_0\ (X\ i)$
shows *stochastic-process* $M\ t_0\ (\lambda k\ \xi.\ \sum_{i \in I}. X\ i\ k\ \xi)$ **using** *assms*[*THEN stochastic-process.random-variable*] **by** (*unfold-locales*, *auto*)

8.1.1 Natural Filtration

The natural filtration induced by a stochastic process X is the filtration generated by all events involving the process up to the time index t , i.e. $\Sigma_t = \sigma\{X_s \mid s \leq t\}$.

definition *natural-filtration* :: *'a measure* \Rightarrow *'b* \Rightarrow (*'b* \Rightarrow *'a* \Rightarrow *'c* :: *topological-space*) \Rightarrow *'b* :: {*second-countable-topology*, *linorder-topology*} \Rightarrow *'a measure*
where

natural-filtration $M\ t_0\ Y = (\lambda t. \text{sigma-gen } (\text{space } M) \text{ borel } \{Y\ i \mid i. i \in \{t_0..t\}\})$

context *stochastic-process*

begin

lemma *sets-natural-filtration'*: *sets* (*natural-filtration* $M\ t_0\ X\ t$) = *sigma-sets* (*space* M) $(\bigcup_{i \in \{t_0..t\}}. \{X\ i - 'A \cap \text{space } M \mid A. A \in \text{borel}\})$
unfolding *natural-filtration-def* *sets-sigma-gen* **by** (*intro sigma-sets-eqI*) *blast+*

lemma

shows *sets-natural-filtration*: *sets* (*natural-filtration* $M\ t_0\ X\ t$) = *sigma-sets* (*space* M) $(\bigcup_{i \in \{t_0..t\}}. \{X\ i - 'A \cap \text{space } M \mid A. \text{open } A\})$

and *space-natural-filtration*[*simp*]: *space* (*natural-filtration* $M\ t_0\ X\ t$) = *space* M

proof –

show *space* (*natural-filtration* $M\ t_0\ X\ t$) = *space* M **unfolding** *natural-filtration-def* *space-sigma-gen* ..

show *sets* (*natural-filtration* $M\ t_0\ X\ t$) = *sigma-sets* (*space* M) $(\bigcup_{i \in \{t_0..t\}}. \{X\ i - 'A \cap \text{space } M \mid A. \text{open } A\})$ **unfolding** *sets-natural-filtration'*

proof (*intro sigma-sets-eqI*, *clarify*)

fix i **and** $A :: 'c \text{ set}$ **assume** *asm*: $i \in \{t_0..t\}$ $A \in \text{sets borel}$

hence $A \in \text{sigma-sets UNIV } \{S. \text{open } S\}$ **unfolding** *borel-def* **by** *simp*

thus $X\ i - 'A \cap \text{space } M \in \text{sigma-sets } (\text{space } M) (\bigcup_{i \in \{t_0..t\}}. \{X\ i - 'A \cap \text{space } M \mid A. \text{open } A\})$

proof (*induction*)

case (*Compl a*)

have $X \text{ } i \text{ } - ' (UNIV - a) \cap \text{space } M = \text{space } M - (X \text{ } i \text{ } - ' a \cap \text{space } M)$ **by** *blast*
then show $?case$ **using** *Compl(2)[THEN sigma-sets.Compl]* **by** *presburger*
next
case (*Union a*)
have $X \text{ } i \text{ } - ' \bigcup (\text{range } a) \cap \text{space } M = \bigcup (\text{range } (\lambda j. X \text{ } i \text{ } - ' a \text{ } j \cap \text{space } M))$
by *blast*
then show $?case$ **using** *Union(2)[THEN sigma-sets.Union]* **by** *presburger*
qed (*auto intro: asm*)
qed (*intro sigma-sets.Basic, fastforce*)
qed

lemma *subalgebra-natural-filtration:*

shows *subalgebra M (natural-filtration M t₀ X i)*
unfolding *subalgebra-def* **using** *measurable-family-iff-contains-sigma-gen* **by** (*force simp add: natural-filtration-def*)

end

sublocale *stochastic-process* \subseteq *filtered-measure-natural-filtration:* *filtered-measure M natural-filtration M t₀ X t₀*

by (*unfold-locales*) (*intro subalgebra-natural-filtration, simp only: sets-natural-filtration, intro sigma-sets-subseteq, force*)

In order to show that the natural filtration constitutes a filtered sigma finite measure, we need to provide a countable exhausting set in the preimage of $X \text{ } t_0$.

lemma (*in sigma-finite-measure*) *sigma-finite-filtered-measure-natural-filtration:*

assumes *stochastic-process M t₀ X*
and *exhausting-set: countable A (bigcup A) = space M \wedge a. a \in A \implies emeasure M a \neq ∞ \wedge a. a \in A \implies $\exists b \in \text{borel. } a = X \text{ } t_0 - ' b \cap \text{space } M$*

shows *sigma-finite-filtered-measure M (natural-filtration M t₀ X) t₀*

proof (*unfold-locales*)

interpret *stochastic-process M t₀ X* **by** (*rule assms*)
have $A \subseteq \text{sets } (\text{restr-to-subalg } M \text{ } (\text{natural-filtration } M \text{ } t_0 \text{ } X \text{ } t_0))$ **using** *exhausting-set* **by** (*simp add: sets-restr-to-subalg[OF subalgebra-natural-filtration] sets-natural-filtration'*)
fast

moreover have $\bigcup A = \text{space } (\text{restr-to-subalg } M \text{ } (\text{natural-filtration } M \text{ } t_0 \text{ } X \text{ } t_0))$

unfolding *space-restr-to-subalg* **using** *exhausting-set* **by** *simp*

moreover have $\forall a \in A. \text{emeasure } (\text{restr-to-subalg } M \text{ } (\text{natural-filtration } M \text{ } t_0 \text{ } X \text{ } t_0)) \text{ } a \neq \infty$ **using** *calculation(1) exhausting-set(3)*

by (*auto simp add: sets-restr-to-subalg[OF subalgebra-natural-filtration] emeasure-restr-to-subalg[OF subalgebra-natural-filtration]*)

ultimately show $\exists A. \text{countable } A \wedge A \subseteq \text{sets } (\text{restr-to-subalg } M \text{ } (\text{natural-filtration } M \text{ } t_0 \text{ } X \text{ } t_0)) \wedge \bigcup A = \text{space } (\text{restr-to-subalg } M \text{ } (\text{natural-filtration } M \text{ } t_0 \text{ } X \text{ } t_0)) \wedge (\forall a \in A. \text{emeasure } (\text{restr-to-subalg } M \text{ } (\text{natural-filtration } M \text{ } t_0 \text{ } X \text{ } t_0)) \text{ } a \neq \infty)$ **using** *exhausting-set* **by** *blast*

show $\bigwedge i j. [t_0 \leq i; i \leq j] \implies \text{sets } (\text{natural-filtration } M \text{ } t_0 \text{ } X \text{ } i) \subseteq \text{sets } (\text{natural-filtration } M \text{ } t_0 \text{ } X \text{ } j)$ **using** *filtered-measure-natural-filtration.subalgebra-F* **by** (*simp add: sub-*

algebra-def)
qed (*auto intro: stochastic-process.subalgebra-natural-filtration assms(1)*)

lemma (*in finite-measure*) *sigma-finite-filtered-measure-natural-filtration*:
 assumes *stochastic-process* $M \ t_0 \ X$
 shows *sigma-finite-filtered-measure* $M \ (natural-filtration \ M \ t_0 \ X) \ t_0$
proof (*intro sigma-finite-filtered-measure-natural-filtration[OF assms(1), of {space M}]*)
 have $space \ M = X \ t_0 - ' UNIV \cap space \ M$ **by** *blast*
 thus $\bigwedge a. a \in \{space \ M\} \implies \exists b \in sets \ borel. a = X \ t_0 - ' b \cap space \ M$ **by** *force*
qed (*auto*)

8.2 Adapted Process

We call a collection a stochastic process X adapted if $X \ i$ is $F \ i$ -borel-measurable for all indices i .

locale *adapted-process* = *filtered-measure* $M \ F \ t_0$ **for** $M \ F \ t_0$ **and** $X :: - \Rightarrow - \Rightarrow -$
 $:: \{second-countable-topology, banach\} +$
 assumes *adapted[measurable]*: $\bigwedge i. t_0 \leq i \implies X \ i \in borel-measurable \ (F \ i)$
begin

lemma *adaptedE[elim]*:
 assumes $\llbracket \bigwedge j \ i. t_0 \leq j \implies j \leq i \implies X \ j \in borel-measurable \ (F \ i) \rrbracket \implies P$
 shows P
using *assms using adapted by (metis dual-order.trans borel-measurable-subalgebra sets-F-mono space-F)*

lemma *adaptedD*:
 assumes $t_0 \leq j \ j \leq i$
 shows $X \ j \in borel-measurable \ (F \ i)$ **using** *assms adaptedE by meson*

end

locale *nat-adapted-process* = *adapted-process* $M \ F \ 0 :: nat \ X$ **for** $M \ F \ X$
sublocale *nat-adapted-process* $\subseteq nat-filtered-measure \ ..$

locale *real-adapted-process* = *adapted-process* $M \ F \ 0 :: real \ X$ **for** $M \ F \ X$
sublocale *real-adapted-process* $\subseteq real-filtered-measure \ ..$

lemma (*in filtered-measure*) *adapted-process-const-fun*:
 assumes $f \in borel-measurable \ (F \ t_0)$
 shows *adapted-process* $M \ F \ t_0 \ (\lambda -. f)$
using *measurable-from-subalg subalgebra-F assms by (unfold-locales) blast*

lemma (*in filtered-measure*) *adapted-process-const*:
 shows *adapted-process* $M \ F \ t_0 \ (\lambda i -. c \ i)$ **by** (*unfold-locales simp*)

context *adapted-process*
begin

lemma *compose*:

assumes $\bigwedge i. t_0 \leq i \implies f\ i \in \text{borel-measurable borel}$
shows *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. (f\ i)\ (X\ i\ \xi))$
by (*unfold-locale*) (*intro measurable-compose*[*OF adapted assms*])

lemma *norm*: *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. \text{norm}\ (X\ i\ \xi))$ **by** (*fastforce intro: compose*)

lemma *scaleR-right*:

assumes *adapted-process* $M\ F\ t_0\ R$
shows *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. (R\ i\ \xi) *_{\mathbb{R}} (X\ i\ \xi))$
using *adapted-process.adapted*[*OF assms*] *adapted by* (*unfold-locale*) *simp*

lemma *scaleR-right-const-fun*:

assumes $f \in \text{borel-measurable}\ (F\ t_0)$
shows *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. f\ \xi *_{\mathbb{R}} (X\ i\ \xi))$
using *assms by* (*fast intro: scaleR-right adapted-process-const-fun*)

lemma *scaleR-right-const*: *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. c\ i *_{\mathbb{R}} (X\ i\ \xi))$ **by** (*unfold-locale*) *simp*

lemma *add*:

assumes *adapted-process* $M\ F\ t_0\ Y$
shows *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi + Y\ i\ \xi)$
using *adapted-process.adapted*[*OF assms*] *adapted by* (*unfold-locale*) *simp*

lemma *diff*:

assumes *adapted-process* $M\ F\ t_0\ Y$
shows *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi - Y\ i\ \xi)$
using *adapted-process.adapted*[*OF assms*] *adapted by* (*unfold-locale*) *simp*

lemma *uminus*: *adapted-process* $M\ F\ t_0\ (-X)$ **using** *scaleR-right-const*[*of* $\lambda -. -1$] **by** (*simp add: fun-Compl-def*)

lemma *partial-sum*: *adapted-process* $M\ F\ t_0\ (\lambda n\ \xi. \sum_{i \in \{t_0..<n\}} X\ i\ \xi)$

proof (*unfold-locale*)

fix $i :: 'b$
have $X\ j \in \text{borel-measurable}\ (F\ i)$ **if** $t_0 \leq j < i$ **for** j **using** *that adaptedE* **by** *fastforce*
thus $(\lambda \xi. \sum_{i \in \{t_0..<i\}} X\ i\ \xi) \in \text{borel-measurable}\ (F\ i)$ **by** *simp*
qed

lemma *partial-sum'*: *adapted-process* $M\ F\ t_0\ (\lambda n\ \xi. \sum_{i \in \{t_0..n\}} X\ i\ \xi)$

proof (*unfold-locale*)

fix $i :: 'b$
have $X\ j \in \text{borel-measurable}\ (F\ i)$ **if** $t_0 \leq j \leq i$ **for** j **using** *that adaptedE* **by** *meson*
thus $(\lambda \xi. \sum_{i \in \{t_0..i\}} X\ i\ \xi) \in \text{borel-measurable}\ (F\ i)$ **by** *simp*

qed

end

lemma (in *nat-adapted-process*) *partial-sum-Suc*: *nat-adapted-process* $M\ F\ (\lambda n\ \xi.$
 $\sum_{i < n}. X\ (Suc\ i)\ \xi)$
proof (*unfold-locales*)
 fix i
 have $X\ j \in \text{borel-measurable}\ (F\ i)$ if $j \leq i$ for j using *that adaptedD* by *blast*
 thus $(\lambda \xi. \sum_{i < i}. X\ (Suc\ i)\ \xi) \in \text{borel-measurable}\ (F\ i)$ by *auto*
 qed

lemma (in *filtered-measure*) *adapted-process-sum*:
 assumes $\bigwedge i. i \in I \implies \text{adapted-process}\ M\ F\ t_0\ (X\ i)$
 shows $\text{adapted-process}\ M\ F\ t_0\ (\lambda k\ \xi. \sum_{i \in I}. X\ i\ k\ \xi)$
proof –
 {
 fix $i\ k$ assume $i \in I$ and *asm*: $t_0 \leq k$
 then interpret $\text{adapted-process}\ M\ F\ t_0\ X\ i$ using *assms* by *simp*
 have $X\ i\ k \in \text{borel-measurable}\ M\ X\ i\ k \in \text{borel-measurable}\ (F\ k)$ using *measurable-from-subalg subalgebra adapted asm* by (*blast, simp*)
 }
 thus ?thesis by (*unfold-locales*) *simp*
 qed

An adapted process is necessarily a stochastic process.

sublocale *adapted-process* \subseteq *stochastic-process* using *measurable-from-subalg subalgebra adapted* by (*unfold-locales*) *blast*

sublocale *nat-adapted-process* \subseteq *nat-stochastic-process* ..

sublocale *real-adapted-process* \subseteq *real-stochastic-process* ..

A stochastic process is always adapted to the natural filtration it generates.

sublocale *stochastic-process* \subseteq *adapted-natural*: *adapted-process* $M\ \text{natural-filtration}\ M\ t_0\ X\ t_0\ X$ by (*unfold-locales*) (*auto simp add: natural-filtration-def intro: random-variable measurable-sigma-gen*)

8.3 Progressively Measurable Process

locale *progressive-process* = *filtered-measure* $M\ F\ t_0$ for $M\ F\ t_0$ and $X :: - \Rightarrow - \Rightarrow - :: \{\text{second-countable-topology, banach}\} +$
 assumes *progressive[measurable]*: $\bigwedge t. t_0 \leq t \implies (\lambda(i, x). X\ i\ x) \in \text{borel-measurable}\ (\text{restrict-space}\ \text{borel}\ \{t_0..t\} \otimes_M F\ t)$
begin

lemma *progressiveD*:
 assumes $S \in \text{borel}$
 shows $(\lambda(j, \xi). X\ j\ \xi) - ' S \cap (\{t_0..i\} \times \text{space}\ M) \in (\text{restrict-space}\ \text{borel}\ \{t_0..i\} \otimes_M F\ i)$

using *measurable-sets*[*OF progressive*, *OF - assms*, *of i*]
by (*cases* $t_0 \leq i$) (*auto simp add: space-F space-restrict-space sets-pair-measure space-pair-measure*)

end

locale *nat-progressive-process* = *progressive-process* $M F 0 :: \text{nat } X$ **for** $M F X$
locale *real-progressive-process* = *progressive-process* $M F 0 :: \text{real } X$ **for** $M F X$

lemma (*in filtered-measure*) *progressive-process-const-fun*:
assumes $f \in \text{borel-measurable } (F t_0)$
shows *progressive-process* $M F t_0 (\lambda \cdot. f)$
proof (*unfold-locales*)
fix i **assume** *asm*: $t_0 \leq i$
have $f \in \text{borel-measurable } (F i)$ **using** *borel-measurable-mono*[*OF order.refl asm*]
assms **by** *blast*
thus *case-prod* $(\lambda \cdot. f) \in \text{borel-measurable } (\text{restrict-space borel } \{t_0..i\} \otimes_M F i)$
using *measurable-compose*[*OF measurable-snd*] **by** *simp*
qed

lemma (*in filtered-measure*) *progressive-process-const*:
assumes $c \in \text{borel-measurable borel}$
shows *progressive-process* $M F t_0 (\lambda i \cdot. c i)$
using *assms* **by** (*unfold-locales*) (*auto simp add: measurable-split-conv intro!:*
measurable-compose[*OF measurable-fst*] *measurable-restrict-space1*)

context *progressive-process*
begin

lemma *compose*:
assumes *case-prod* $f \in \text{borel-measurable borel}$
shows *progressive-process* $M F t_0 (\lambda i \xi. (f i) (X i \xi))$
proof
fix i **assume** *asm*: $t_0 \leq i$
have $(\lambda(j, \xi). (j, X j \xi)) \in (\text{restrict-space borel } \{t_0..i\} \otimes_M F i) \rightarrow_M \text{borel } \otimes_M \text{borel}$
using *progressive*[*OF asm*] *measurable-fst''*[*OF measurable-restrict-space1*, *OF measurable-id*]
by (*auto simp add: measurable-pair-iff measurable-split-conv*)
moreover **have** $(\lambda(j, \xi). f j (X j \xi)) = \text{case-prod } f \circ ((\lambda(j, y). (j, y)) \circ (\lambda(j, \xi). (j, X j \xi)))$ **by** *fastforce*
ultimately show $(\lambda(j, \xi). (f j) (X j \xi)) \in \text{borel-measurable } (\text{restrict-space borel } \{t_0..i\} \otimes_M F i)$ **using** *assms* **by** (*simp add: borel-prod*)
qed

lemma *norm*: *progressive-process* $M F t_0 (\lambda i \xi. \text{norm } (X i \xi))$ **using** *measurable-compose*[*OF progressive borel-measurable-norm*] **by** (*unfold-locales*) *simp*

lemma *scaleR-right*:

assumes *progressive-process* $M F t_0 R$
shows *progressive-process* $M F t_0 (\lambda i \xi. (R i \xi) *_R (X i \xi))$
using *progressive-process.progressive*[*OF assms*] **by** (*unfold-locales*) (*simp add: progressive assms*)

lemma *scaleR-right-const-fun*:
assumes $f \in \text{borel-measurable } (F t_0)$
shows *progressive-process* $M F t_0 (\lambda i \xi. f \xi *_R (X i \xi))$
using *assms* **by** (*fastforce intro: scaleR-right progressive-process-const-fun*)

lemma *scaleR-right-const*:
assumes $c \in \text{borel-measurable borel}$
shows *progressive-process* $M F t_0 (\lambda i \xi. c i *_R (X i \xi))$
using *assms* **by** (*fastforce intro: scaleR-right progressive-process-const*)

lemma *add*:
assumes *progressive-process* $M F t_0 Y$
shows *progressive-process* $M F t_0 (\lambda i \xi. X i \xi + Y i \xi)$
using *progressive-process.progressive*[*OF assms*] **by** (*unfold-locales*) (*simp add: progressive assms*)

lemma *diff*:
assumes *progressive-process* $M F t_0 Y$
shows *progressive-process* $M F t_0 (\lambda i \xi. X i \xi - Y i \xi)$
using *progressive-process.progressive*[*OF assms*] **by** (*unfold-locales*) (*simp add: progressive assms*)

lemma *uminus: progressive-process* $M F t_0 (-X)$ **using** *scaleR-right-const*[*of λ-. -1*] **by** (*simp add: fun-Compl-def*)

end

A progressively measurable process is also adapted.

sublocale *progressive-process* \subseteq *adapted-process* **using** *measurable-compose-rev*[*OF progressive measurable-Pair1*] **unfolding** *prod.case* **by** (*unfold-locales*) *simp*

sublocale *nat-progressive-process* \subseteq *nat-adapted-process* ..

sublocale *real-progressive-process* \subseteq *real-adapted-process* ..

In the discrete setting, adaptedness is equivalent to progressive measurability.

sublocale *nat-adapted-process* \subseteq *nat-progressive-process*

proof (*unfold-locales, intro borel-measurableI*)

fix $S :: 'b \text{ set}$ **and** $i :: \text{nat}$ **assume** *open-S: open S*
{
fix j **assume** *asm: $j \leq i$*
hence $X j - 'S \cap \text{space } M \in F i$ **using** *adaptedD*[*of j, THEN measurable-sets*]
space-F open-S **by** *fastforce*

moreover have $\text{case-prod } X - 'S \cap \{j\} \times \text{space } M = \{j\} \times (X j - 'S \cap \text{space } M)$ **for** j **by** *fast*
moreover have $\{j :: \text{nat}\} \in \text{restrict-space borel } \{0..i\}$ **using** *asm* **by** (*simp* *add: sets-restrict-space-iff*)
ultimately have $\text{case-prod } X - 'S \cap \{j\} \times \text{space } M \in \text{restrict-space borel } \{0..i\} \otimes_M F i$ **by** *simp*
}
hence $(\lambda j. (\lambda(x, y). X x y) - 'S \cap \{j\} \times \text{space } M) ' \{..i\} \subseteq \text{restrict-space borel } \{0..i\} \otimes_M F i$ **by** *blast*
moreover have $\text{case-prod } X - 'S \cap \text{space } (\text{restrict-space borel } \{0..i\} \otimes_M F i) = (\bigcup_{j \leq i}. \text{case-prod } X - 'S \cap \{j\} \times \text{space } M)$ **unfolding** *space-pair-measure space-restrict-space space-F* **by** *force*
ultimately show $\text{case-prod } X - 'S \cap \text{space } (\text{restrict-space borel } \{0..i\} \otimes_M F i) \in \text{restrict-space borel } \{0..i\} \otimes_M F i$ **by** (*metis sets.countable-UN*)
qed

8.4 Predictable Process

We introduce the constant Σ_P to denote the predictable sigma algebra.

context *filtered-measure*
begin

definition $\Sigma_P :: ('b \times 'a) \text{ measure where predictable-sigma: } \Sigma_P \equiv \text{sigma } (\{t_0..\} \times \text{space } M) (\{\{s <..t\} \times A \mid A s t. A \in F s \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in F t_0\})$

lemma *space-predictable-sigma[simp]: space $\Sigma_P = (\{t_0..\} \times \text{space } M)$ unfolding predictable-sigma space-measure-of-conv* **by** *blast*

lemma *sets-predictable-sigma: sets $\Sigma_P = \text{sigma-sets } (\{t_0..\} \times \text{space } M) (\{\{s <..t\} \times A \mid A s t. A \in F s \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in F t_0\})$*

unfolding *predictable-sigma using space-F sets.sets-into-space* **by** (*subst sets-measure-of fastforce+*)

lemma *measurable-predictable-sigma-snd:*

assumes *countable \mathcal{I} $\mathcal{I} \subseteq \{\{s <..t\} \mid s t. t_0 \leq s \wedge s < t\} \{t_0 <..\} \subseteq (\bigcup \mathcal{I})$*

shows $\text{snd} \in \Sigma_P \rightarrow_M F t_0$

proof (*intro measurableI, force simp add: space-F*)

fix $S :: 'a \text{ set assume } \text{asm: } S \in F t_0$

have *countable: countable $(\lambda I. I \times S) ' \mathcal{I}$* **using** *assms(1)* **by** *blast*

have $(\lambda I. I \times S) ' \mathcal{I} \subseteq \{\{s <..t\} \times A \mid A s t. A \in F s \wedge t_0 \leq s \wedge s < t\}$ **using** *sets-F-mono[OF order-refl, THEN subsetD, OF - asm]* *assms(2)* **by** *blast*

hence $(\bigcup I \in \mathcal{I}. I \times S) \cup \{t_0\} \times S \in \Sigma_P$ **unfolding** *sets-predictable-sigma* **using** *asm* **by** (*intro sigma-sets-Un[OF sigma-sets-UNION[OF countable] sigma-sets.Basic] sigma-sets.Basic*) *blast+*

moreover have $\text{snd} - 'S \cap \text{space } \Sigma_P = \{t_0..\} \times S$ **using** *sets.sets-into-space[OF asm]* **by** (*fastforce simp add: space-F*)

moreover have $(\bigcup I \in \mathcal{I}. I \times S) \cup \{t_0\} \times S = \{t_0..\} \times S$ **using** *assms(2,3)* **using** *ivl-disj-un(1)* **by** *fastforce*

ultimately show $snd - ' S \cap space \Sigma_P \in \Sigma_P$ by argo
qed

lemma *measurable-predictable-sigma-fst*:

assumes *countable* \mathcal{I} $\mathcal{I} \subseteq \{\{s<..t\} \mid s \ t. \ t_0 \leq s \wedge s < t\} \ \{t_0<..\} \subseteq (\bigcup \mathcal{I})$

shows $fst \in \Sigma_P \rightarrow_M \text{borel}$

proof –

have $A \times space \ M \in sets \ \Sigma_P$ if $A \in sigma\text{-sets} \ \{t_0..\} \ \{\{s<..t\} \mid s \ t. \ t_0 \leq s \wedge s < t\}$ for A **unfolding** *sets-predictable-sigma* **using** that

proof (*induction rule: sigma-sets.induct*)

case (*Basic a*)

thus ?case **using** *space-F sets.top* **by** *blast*

next

case (*Compl a*)

have $(\{t_0..\} - a) \times space \ M = \{t_0..\} \times space \ M - a \times space \ M$ **by** *blast*

then show ?case **using** *Compl(2)* [*THEN sigma-sets.Compl*] **by** *presburger*

next

case (*Union a*)

have $\bigcup (range \ a) \times space \ M = \bigcup (range \ (\lambda i. \ a \ i \times space \ M))$ **by** *blast*

then show ?case **using** *Union(2)* [*THEN sigma-sets.Union*] **by** *presburger*

qed (*auto*)

moreover have *restrict-space borel* $\{t_0..\} = sigma \ \{t_0..\} \ \{\{s<..t\} \mid s \ t. \ t_0 \leq s \wedge s < t\}$

proof –

have *sigma-sets* $\{t_0..\} \ ((\cap) \ \{t_0..\} \ ' sigma\text{-sets} \ UNIV \ (range \ greaterThan)) = sigma\text{-sets} \ \{t_0..\} \ \{\{s<..t\} \mid s \ t. \ t_0 \leq s \wedge s < t\}$

proof (*intro sigma-sets-eqI ; clarify*)

fix $A :: 'b \ set$ **assume** *asm*: $A \in sigma\text{-sets} \ UNIV \ (range \ greaterThan)$

thus $\{t_0..\} \cap A \in sigma\text{-sets} \ \{t_0..\} \ \{\{s<..t\} \mid s \ t. \ t_0 \leq s \wedge s < t\}$

proof (*induction rule: sigma-sets.induct*)

case (*Basic a*)

then obtain s **where** $a = \{s<..\}$ **by** *blast*

show ?case

proof (*cases* $t_0 \leq s$)

case *True*

hence $*$: $\{t_0..\} \cap a = (\bigcup i \in \mathcal{I}. \ \{s<..\} \cap i)$ **using** *s assms(3)* **by** *force*

have $((\cap) \ \{s<..\} \ ' \mathcal{I}) \subseteq sigma\text{-sets} \ \{t_0..\} \ \{\{s<..t\} \mid s \ t. \ t_0 \leq s \wedge s < t\}$

proof (*clarify*)

fix A **assume** $A \in \mathcal{I}$

then obtain $s' \ t'$ **where** $A = \{s'<..t'\} \ t_0 \leq s' \ s' < t'$ **using** *assms(2)*

by *blast*

hence $\{s<..\} \cap A = \{\max s \ s'<..t'\}$ **by** *fastforce*

moreover have $t_0 \leq \max s \ s'$ **using** *A True* **by** *linarith*

moreover have $\max s \ s' < t'$ **if** $s < t'$ **using** *A that* **by** *linarith*

moreover have $\{s<..\} \cap A = \{\}$ **if** $\neg s < t'$ **using** *A that* **by** *force*

ultimately show $\{s<..\} \cap A \in sigma\text{-sets} \ \{t_0..\} \ \{\{s<..t\} \mid s \ t. \ t_0 \leq s \wedge s < t\}$ **by** (*cases* $s < t'$) (*blast, simp add: sigma-sets.Empty*)

qed

thus ?thesis **unfolding** $*$ **using** *assms(1)* **by** (*intro sigma-sets-UNION*)

```

auto
  next
    case False
    hence  $\{t_0..\} \cap a = \{t_0..\}$  using s by force
    thus ?thesis using sigma-sets-top by auto
  qed
  next
    case (Compl a)
    have  $\{t_0..\} \cap (UNIV - a) = \{t_0..\} - (\{t_0..\} \cap a)$  by blast
    then show ?case using Compl(2)[THEN sigma-sets.Compl] by presburger
  next
    case (Union a)
    have  $\{t_0..\} \cap \bigcup (\text{range } a) = \bigcup (\text{range } (\lambda i. \{t_0..\} \cap a \ i))$  by blast
    then show ?case using Union(2)[THEN sigma-sets.Union] by presburger
  qed (simp add: sigma-sets.Empty)
  next
    fix s t assume asm:  $t_0 \leq s < t$ 
    hence *:  $\{s<..t\} = \{s<..\} \cap (\{t_0..\} - \{t<..\})$  by force
    have  $\{s<..\} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\})$  ' sigma-sets UNIV (range greaterThan))
  using asm by (intro sigma-sets.Basic) auto
    moreover have  $\{t_0..\} - \{t<..\} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\})$  ' sigma-sets UNIV (range greaterThan))
  using asm by (intro sigma-sets.Compl sigma-sets.Basic) auto
  auto
    ultimately show  $\{s<..t\} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\})$  ' sigma-sets UNIV (range greaterThan))
  unfolding * Int-range-binary[of  $\{s<..\}$ ] by (intro sigma-sets.Inter[OF - binary-in-sigma-sets]) auto
  qed
    thus ?thesis unfolding borel-Ioi restrict-space-def emeasure-sigma by (force intro: sigma-eqI)
  qed
    ultimately have  $\text{restrict-space borel } \{t_0..\} \otimes_M \text{sigma } (\text{space } M) \{ \} \subseteq \text{sets } \Sigma_P$ 

    unfolding sets-pair-measure space-restrict-space space-measure-of-conv
    using space-predictable-sigma sets.sigma-algebra-axioms[of  $\Sigma_P$ ]
    by (intro sigma-algebra.sigma-sets-subset) (auto simp add: sigma-sets-empty-eq sets-measure-of-conv)
    moreover have  $\text{space } (\text{restrict-space borel } \{t_0..\} \otimes_M \text{sigma } (\text{space } M) \{ \}) = \text{space } \Sigma_P$ 
  by (simp add: space-pair-measure)
    moreover have  $\text{fst} \in \text{restrict-space borel } \{t_0..\} \otimes_M \text{sigma } (\text{space } M) \{ \} \rightarrow_M \text{borel}$ 
  by (fastforce intro: measurable-fst'[OF measurable-restrict-space1, of  $\lambda x. x$ ])

    ultimately show ?thesis by (meson borel-measurable-subalgebra)
  qed
end

locale predictable-process = filtered-measure M F t_0 for M F t_0 and X ::  $- \Rightarrow - \Rightarrow - :: \{ \text{second-countable-topology, banach} \} +$ 
  assumes predictable:  $(\lambda(t, x). X \ t \ x) \in \text{borel-measurable } \Sigma_P$ 

```

begin

lemmas *predictableD = measurable-sets[OF predictable, unfolded space-predictable-sigma]*

end

locale *nat-predictable-process = predictable-process M F 0 :: nat X for M F X*

locale *real-predictable-process = predictable-process M F 0 :: real X for M F X*

lemma (**in** *nat-filtered-measure*) *measurable-predictable-sigma-snd:*

shows *snd* $\in \Sigma_P \rightarrow_M F 0$

by (*intro measurable-predictable-sigma-snd[of range ($\lambda x. \{Suc\ x\}$)]*) (*force | simp add: greaterThan-0*)**+**

lemma (**in** *nat-filtered-measure*) *measurable-predictable-sigma-fst:*

shows *fst* $\in \Sigma_P \rightarrow_M \text{borel}$

by (*intro measurable-predictable-sigma-fst[of range ($\lambda x. \{Suc\ x\}$)]*) (*force | simp add: greaterThan-0*)**+**

lemma (**in** *real-filtered-measure*) *measurable-predictable-sigma-snd:*

shows *snd* $\in \Sigma_P \rightarrow_M F 0$

using *real-arch-simple* **by** (*intro measurable-predictable-sigma-snd[of range ($\lambda x::nat. \{0 <..real\ (Suc\ x)\}$)]*) (*fastforce intro: add-increasing*)**+**

lemma (**in** *real-filtered-measure*) *measurable-predictable-sigma-fst:*

shows *fst* $\in \Sigma_P \rightarrow_M \text{borel}$

using *real-arch-simple* **by** (*intro measurable-predictable-sigma-fst[of range ($\lambda x::nat. \{0 <..real\ (Suc\ x)\}$)]*) (*fastforce intro: add-increasing*)**+**

lemma (**in** *filtered-measure*) *predictable-process-const-fun:*

assumes *snd* $\in \Sigma_P \rightarrow_M F t_0$ *f* $\in \text{borel-measurable}\ (F\ t_0)$

shows *predictable-process M F t₀* ($\lambda\cdot. f$)

using *measurable-compose-rev[OF assms(2)] assms(1)* **by** (*unfold-locales*) (*auto simp add: measurable-split-conv*)

lemma (**in** *nat-filtered-measure*) *predictable-process-const-fun:*

assumes *f* $\in \text{borel-measurable}\ (F\ 0)$

shows *nat-predictable-process M F* ($\lambda\cdot. f$)

using *assms* **by** (*intro predictable-process-const-fun[OF measurable-predictable-sigma-snd, THEN nat-predictable-process.intro]*)

lemma (**in** *real-filtered-measure*) *predictable-process-const-fun:*

assumes *f* $\in \text{borel-measurable}\ (F\ 0)$

shows *real-predictable-process M F* ($\lambda\cdot. f$)

using *assms* **by** (*intro predictable-process-const-fun[OF measurable-predictable-sigma-snd,*

THEN real-predictable-process.intro)

lemma (in *filtered-measure*) *predictable-process-const*:
 assumes $fst \in \text{borel-measurable } \Sigma_P \ c \in \text{borel-measurable borel}$
 shows *predictable-process* $M \ F \ t_0 \ (\lambda i \ -. \ c \ i)$
 using *assms* by (unfold-locale) (simp add: measurable-split-conv)

lemma (in *filtered-measure*) *predictable-process-const'*:
 shows *predictable-process* $M \ F \ t_0 \ (\lambda \cdot \ -. \ c)$
 by (unfold-locale) simp

lemma (in *nat-filtered-measure*) *predictable-process-const*:
 assumes $c \in \text{borel-measurable borel}$
 shows *nat-predictable-process* $M \ F \ (\lambda i \ -. \ c \ i)$
 using *assms* by (intro *predictable-process-const*[OF *measurable-predictable-sigma-fst*,
THEN nat-predictable-process.intro])

lemma (in *real-filtered-measure*) *predictable-process-const*:
 assumes $c \in \text{borel-measurable borel}$
 shows *real-predictable-process* $M \ F \ (\lambda i \ -. \ c \ i)$
 using *assms* by (intro *predictable-process-const*[OF *measurable-predictable-sigma-fst*,
THEN real-predictable-process.intro])

context *predictable-process*
begin

lemma *compose*:
 assumes $fst \in \text{borel-measurable } \Sigma_P \ \text{case-prod } f \in \text{borel-measurable borel}$
 shows *predictable-process* $M \ F \ t_0 \ (\lambda i \ \xi. \ (f \ i) \ (X \ i \ \xi))$
proof
 have $(\lambda(i, \xi). \ (i, X \ i \ \xi)) \in \Sigma_P \rightarrow_M \text{borel} \otimes_M \text{borel}$ using *predictable assms*(1)
 by (auto simp add: measurable-pair-iff measurable-split-conv)
 moreover have $(\lambda(i, \xi). \ f \ i \ (X \ i \ \xi)) = \text{case-prod } f \ o \ (\lambda(i, \xi). \ (i, X \ i \ \xi))$ by
fastforce
 ultimately show $(\lambda(i, \xi). \ f \ i \ (X \ i \ \xi)) \in \text{borel-measurable } \Sigma_P$ unfolding *borel-prod*
 using *assms* by simp
qed

lemma *norm*: *predictable-process* $M \ F \ t_0 \ (\lambda i \ \xi. \ \text{norm} \ (X \ i \ \xi))$ using *measurable-compose*[OF *predictable borel-measurable-norm*]
 by (unfold-locale) (simp add: prod.case-distrib)

lemma *scaleR-right*:
 assumes *predictable-process* $M \ F \ t_0 \ R$
 shows *predictable-process* $M \ F \ t_0 \ (\lambda i \ \xi. \ (R \ i \ \xi) *_{\mathbb{R}} (X \ i \ \xi))$
 using *predictable predictable-process.predictable*[OF *assms*] by (unfold-locale)
 (auto simp add: measurable-split-conv)

lemma *scaleR-right-const-fun*:

assumes $snd \in \Sigma_P \rightarrow_M F\ t_0\ f \in \text{borel-measurable}\ (F\ t_0)$
shows $\text{predictable-process}\ M\ F\ t_0\ (\lambda i\ \xi. f\ \xi *_{\mathcal{R}} (X\ i\ \xi))$
using *assms* **by** (*fast intro: scaleR-right predictable-process-const-fun*)

lemma *scaleR-right-const*:
assumes $fst \in \text{borel-measurable}\ \Sigma_P\ c \in \text{borel-measurable}\ \text{borel}$
shows $\text{predictable-process}\ M\ F\ t_0\ (\lambda i\ \xi. c\ i *_{\mathcal{R}} (X\ i\ \xi))$
using *assms* **by** (*fastforce intro: scaleR-right predictable-process-const*)

lemma *scaleR-right-const'*: $\text{predictable-process}\ M\ F\ t_0\ (\lambda i\ \xi. c *_{\mathcal{R}} (X\ i\ \xi))$
by (*fastforce intro: scaleR-right predictable-process-const'*)

lemma *add*:
assumes $\text{predictable-process}\ M\ F\ t_0\ Y$
shows $\text{predictable-process}\ M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi + Y\ i\ \xi)$
using *predictable predictable-process.predictable[OF assms]* **by** (*unfold-locales*)
(auto simp add: measurable-split-conv)

lemma *diff*:
assumes $\text{predictable-process}\ M\ F\ t_0\ Y$
shows $\text{predictable-process}\ M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi - Y\ i\ \xi)$
using *predictable predictable-process.predictable[OF assms]* **by** (*unfold-locales*)
(auto simp add: measurable-split-conv)

lemma *uminus*: $\text{predictable-process}\ M\ F\ t_0\ (-X)$ **using** *scaleR-right-const'[of -1]*
by (*simp add: fun-Compl-def*)

end

Every predictable process is also progressively measurable.

sublocale $\text{predictable-process} \subseteq \text{progressive-process}$

proof (*unfold-locales*)

fix $i :: 'b$ **assume** $asm: t_0 \leq i$
 $\{$
fix $S :: ('b \times 'a)$ *set* **assume** $S \in \{\{s <..t\} \times A \mid A\ s\ t. A \in F\ s \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in F\ t_0\}$
hence $(\lambda x. x) - ' S \cap (\{t_0..i\} \times \text{space}\ M) \in \text{restrict-space borel}\ \{t_0..i\} \otimes_M F$
 i

proof

assume $S \in \{\{s <..t\} \times A \mid A\ s\ t. A \in F\ s \wedge t_0 \leq s \wedge s < t\}$
then obtain $s\ t\ A$ **where** $S\text{-is: } S = \{s <..t\} \times A\ t_0 \leq s\ s < t\ A \in F\ s$ **by**
blast

hence $(\lambda x. x) - ' S \cap (\{t_0..i\} \times \text{space}\ M) = \{s <.. \min\ i\ t\} \times A$ **using**
sets.sets-into-space[OF S-is(4)] **by** (*auto simp add: space-F*)

then show $?thesis$ **using** $S\text{-is sets-F-mono[of s i]$ **by** (*cases s ≤ i*) (*fastforce simp add: sets-restrict-space-iff*)+

next

assume $S \in \{\{t_0\} \times A \mid A. A \in F\ t_0\}$
then obtain A **where** $S\text{-is: } S = \{t_0\} \times A\ A \in F\ t_0$ **by** *blast*

hence $(\lambda x. x) - ' S \cap (\{t_0..i\} \times \text{space } M) = \{t_0\} \times A$ **using** *asm sets.sets-into-space*[*OF S-is(2)*] **by** (*auto simp add: space-F*)
thus *?thesis using S-is(2) sets-F-mono*[*OF order-refl asm*] *asm by* (*fastforce simp add: sets-restrict-space-iff*)
qed
hence $(\lambda x. x) - ' S \cap \text{space } (\text{restrict-space borel } \{t_0..i\} \otimes_M F i) \in \text{restrict-space borel } \{t_0..i\} \otimes_M F i$ **by** (*simp add: space-pair-measure space-F*[*OF asm*])
}
moreover have $\{\{s<..t\} \times A \mid A s t. A \in \text{sets } (F s) \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in \text{sets } (F t_0)\} \subseteq \text{Pow } (\{t_0..i\} \times \text{space } M)$ **using** *sets.sets-into-space by* (*fastforce simp add: space-F*)
ultimately have $(\lambda x. x) \in \text{restrict-space borel } \{t_0..i\} \otimes_M F i \rightarrow_M \Sigma_P$ **using** *space-F*[*OF asm*] **by** (*intro measurable-sigma-sets*[*OF sets-predictable-sigma*]) (*fast, force simp add: space-pair-measure*)
thus *case-prod* $X \in \text{borel-measurable } (\text{restrict-space borel } \{t_0..i\} \otimes_M F i)$ **using** *predictable by simp*
qed

sublocale *nat-predictable-process* \subseteq *nat-progressive-process* ..
sublocale *real-predictable-process* \subseteq *real-progressive-process* ..

The following lemma characterizes predictability in a discrete-time setting.

lemma (*in nat-filtered-measure*) *sets-in-filtration*:

assumes $(\bigcup i. \{i\} \times A i) \in \Sigma_P$
shows $A (\text{Suc } i) \in F i \wedge 0 \in F 0$
using *assms unfolding sets-predictable-sigma*
proof (*induction* $(\bigcup i. \{i\} \times A i)$ *arbitrary: A*)
case Basic
{
assume $\exists S. (\bigcup i. \{i\} \times A i) = \{0\} \times S$
then obtain S **where** $S: (\bigcup i. \{i\} \times A i) = \{bot\} \times S$ **unfolding** *bot-nat-def*
by *blast*
hence $S \in F bot$ **using** *Basic by* (*fastforce simp add: times-eq-iff bot-nat-def*)
moreover have $A i = \{\}$ **if** $i \neq bot$ **for** i **using** *that S by blast*
moreover have $A bot = S$ **using** *S by blast*
ultimately have $A (\text{Suc } i) \in F i \wedge 0 \in F 0$ **for** i **unfolding** *bot-nat-def by* (*auto simp add: bot-nat-def*)
}
note $*$ **=** *this*
{
assume $\nexists S. (\bigcup i. \{i\} \times A i) = \{0\} \times S$
then obtain $s t B$ **where** $B: (\bigcup i. \{i\} \times A i) = \{s<..t\} \times B$ $B \in \text{sets } (F s)$ $s < t$ **using** *Basic by auto*
hence $A i = B$ **if** $i \in \{s<..t\}$ **for** i **using** *that by fast*
moreover have $A i = \{\}$ **if** $i \notin \{s<..t\}$ **for** i **using** *B that by fastforce*
ultimately have $A (\text{Suc } i) \in F i \wedge 0 \in F 0$ **for** i **unfolding** *bot-nat-def using* *B sets-F-mono by* (*auto simp add: bot-nat-def*) (*metis less-Suc-eq-le sets.empty-sets subset-eq*)
}


```

note ** = this
show  $A (Suc\ i) \in sets\ (F\ i)\ A\ 0 \in sets\ (F\ 0)$  using  $*(1)[of\ i]\ *(2)\ ***(1)[of\ i]$ 
 $**(2)$  by blast+
next
  case Empty
  {
    case 1
    then show ?case using Empty by simp
  next
    case 2
    then show ?case using Empty by simp
  }
next
  case (Compl a)
  have a-in:  $a \subseteq \{0..\} \times space\ M$  using Compl(1) sets.sets-into-space sets-predictable-sigma
space-predictable-sigma by metis
  hence A-in:  $A\ i \subseteq space\ M$  for i using Compl(4) by blast
  have a:  $a = \{0..\} \times space\ M - (\bigcup i. \{i\} \times A\ i)$  using a-in Compl(4) by blast
  also have  $\dots = - (\bigcap j. - (\{j\} \times (space\ M - A\ j)))$  by blast
  also have  $\dots = (\bigcup j. \{j\} \times (space\ M - A\ j))$  by blast
  finally have  $*$ :  $(space\ M - A\ (Suc\ i)) \in F\ i\ (space\ M - A\ 0) \in F\ 0$  using
Compl(2,3) by auto
  {
    case 1
    then show ?case using  $*$  A-in by (metis bot-nat-0.extremum double-diff
sets.Diff sets.top sets-F-mono sets-le-imp-space-le space-F)
  next
    case 2
    then show ?case using  $*$  A-in by (metis bot-nat-0.extremum double-diff
sets.Diff sets.top sets-F-mono sets-le-imp-space-le space-F)
  }
next
  case (Union a)
  have a-in:  $a\ i \subseteq \{0..\} \times space\ M$  for i using Union(1) sets.sets-into-space
sets-predictable-sigma space-predictable-sigma by metis
  hence A-in:  $A\ i \subseteq space\ M$  for i using Union(4) by blast
  have snd x  $\in snd\ ' (a\ i \cap (\{fst\ x\} \times space\ M))$  if  $x \in a\ i$  for i x using that
a-in by fastforce
  hence a-i:  $a\ i = (\bigcup j. \{j\} \times (snd\ ' (a\ i \cap (\{j\} \times space\ M))))$  for i by force
  have A-i:  $A\ i = snd\ ' (\bigcup (range\ a) \cap (\{i\} \times space\ M))$  for i unfolding
Union(4) using A-in by force
  have  $*$ :  $snd\ ' (a\ j \cap (\{Suc\ i\} \times space\ M)) \in F\ i\ snd\ ' (a\ j \cap (\{0\} \times space\ M))$ 
 $\in F\ 0$  for j using Union(2,3)[OF a-i] by auto
  {
    case 1
    have  $(\bigcup j. snd\ ' (a\ j \cap (\{Suc\ i\} \times space\ M))) \in F\ i$  using  $*$  by fast
    moreover have  $(\bigcup j. snd\ ' (a\ j \cap (\{Suc\ i\} \times space\ M))) = snd\ ' (\bigcup (range\ a) \cap (\{Suc\ i\} \times space\ M))$  by fast
    ultimately show ?case using A-i by metis
  }

```

```

next
  case 2
  have ( $\bigcup j. \text{snd } (a \ j \cap (\{0\} \times \text{space } M)) \in F \ 0$ ) using * by fast
  moreover have ( $\bigcup j. \text{snd } (a \ j \cap (\{0\} \times \text{space } M)) = \text{snd } ( \bigcup (\text{range } a) \cap (\{0\} \times \text{space } M) )$ ) by fast
  ultimately show ?case using A-i by metis
}
qed

```

This leads to the following useful fact.

```

lemma (in nat-predictable-process) adapted-Suc: nat-adapted-process  $M \ F \ (\lambda i. X \ (Suc \ i))$ 
proof (unfold-locales, intro borel-measurableI)
  fix  $S :: 'b \ set$  and  $i$  assume open-S: open  $S$ 
  have  $\{Suc \ i\} = \{i <.. Suc \ i\}$  by fastforce
  hence  $\{Suc \ i\} \times \text{space } M \in \Sigma_P$  unfolding space-F[symmetric, of  $i$ ] sets-predictable-sigma
by (intro sigma-sets.Basic) blast
  moreover have case-prod  $X - ' S \cap (UNIV \times \text{space } M) \in \Sigma_P$  unfolding
atLeast-0[symmetric] using open-S by (intro predictableD, simp add: borel-open)
  ultimately have case-prod  $X - ' S \cap (\{Suc \ i\} \times \text{space } M) \in \Sigma_P$  unfolding
sets-predictable-sigma using space-F sets.sets-into-space
  by (subst Times-Int-distrib1[of  $\{Suc \ i\}$  UNIV space  $M$ , simplified], subst
inf commute, subst Int-assoc[symmetric], subst Int-range-binary)
  (intro sigma-sets-Inter binary-in-sigma-sets, fast)+
  moreover have case-prod  $X - ' S \cap (\{Suc \ i\} \times \text{space } M) = \{Suc \ i\} \times (X \ (Suc \ i) - ' S \cap \text{space } M)$  by (auto simp add: le-Suc-eq)
  moreover have ... = ( $\bigcup j. \{j\} \times (\text{if } j = Suc \ i \text{ then } (X \ (Suc \ i) - ' S \cap \text{space } M) \text{ else } \{\})$ ) by (force split: if-splits)
  ultimately have ( $\bigcup j. \{j\} \times (\text{if } j = Suc \ i \text{ then } (X \ (Suc \ i) - ' S \cap \text{space } M) \text{ else } \{\})$ )  $\in \Sigma_P$  by argo
  thus  $X \ (Suc \ i) - ' S \cap \text{space } (F \ i) \in \text{sets } (F \ i)$  using sets-in-filtration[of  $\lambda j. \text{if } j = Suc \ i \text{ then } (X \ (Suc \ i) - ' S \cap \text{space } M) \text{ else } \{\}$ ] space-F] by presburger
qed

```

theorem nat-predictable-process-iff: nat-predictable-process $M \ F \ X \longleftrightarrow$ nat-adapted-process $M \ F \ (\lambda i. X \ (Suc \ i)) \wedge X \ 0 \in \text{borel-measurable } (F \ 0)$

```

proof (intro iffI)
  assume asm: nat-adapted-process  $M \ F \ (\lambda i. X \ (Suc \ i)) \wedge X \ 0 \in \text{borel-measurable } (F \ 0)$ 
  interpret nat-adapted-process  $M \ F \ \lambda i. X \ (Suc \ i)$  using asm by blast
  have  $(\lambda(x, y). X \ x \ y) \in \text{borel-measurable } \Sigma_P$ 
  proof (intro borel-measurableI)
    fix  $S :: 'b \ set$  assume open-S: open  $S$ 
    have  $\{i\} \times (X \ i - ' S \cap \text{space } M) \in \text{sets } \Sigma_P$  for  $i$ 
    proof (cases  $i$ )
      case 0
      then show ?thesis unfolding sets-predictable-sigma
      using measurable-sets[OF - borel-open[OF open-S], of  $X \ 0 \ F \ 0$ ] asm
      by (auto simp add: space-F)

```

```

next
  case (Suc i)
  have {Suc i} = {i<..Suc i} by fastforce
  then show ?thesis unfolding sets-predictable-sigma
    using measurable-sets[OF adapted borel-open[OF open-S], of i]
    by (intro sigma-sets.Basic, auto simp add: space-F Suc)
qed
moreover have ( $\lambda(x, y). X\ x\ y$ ) - '  $S \cap \text{space } \Sigma_P = (\bigcup i. \{i\} \times (X\ i - ' S \cap \text{space } M))$  by fastforce
ultimately show ( $\lambda(x, y). X\ x\ y$ ) - '  $S \cap \text{space } \Sigma_P \in \text{sets } \Sigma_P$  by simp
qed
thus nat-predictable-process M F X by (unfold-locales)
next
  assume asm: nat-predictable-process M F X
  interpret nat-predictable-process M F X by (rule asm)
  show nat-adapted-process M F ( $\lambda i. X\ (Suc\ i)$ )  $\wedge X\ 0 \in \text{borel-measurable } (F\ 0)$ 
using adapted-Suc by simp
qed

end
theory Martingale
  imports Stochastic-Process Conditional-Expectation-Banach
begin

```

9 Martingales

The following locales are necessary for defining martingales.

locale *sigma-finite-adapted-process* = *adapted-process* + *sigma-finite-filtered-measure*

locale *nat-sigma-finite-adapted-process* = *sigma-finite-adapted-process* M F 0 :: nat X for M F X

locale *real-sigma-finite-adapted-process* = *sigma-finite-adapted-process* M F 0 :: real X for M F X

sublocale *nat-sigma-finite-adapted-process* \subseteq *nat-sigma-finite-filtered-measure* ..

sublocale *real-sigma-finite-adapted-process* \subseteq *real-sigma-finite-filtered-measure* ..

locale *sigma-finite-adapted-process-order* = *sigma-finite-adapted-process* M F t_0 X for M F t_0 and X :: - \Rightarrow - \Rightarrow - :: {order-topology, ordered-real-vector}

locale *nat-sigma-finite-adapted-process-order* = *sigma-finite-adapted-process-order* M F 0 :: nat X for M F X

locale *real-sigma-finite-adapted-process-order* = *sigma-finite-adapted-process-order* M F 0 :: real X for M F X

sublocale *nat-sigma-finite-adapted-process-order* \subseteq *nat-sigma-finite-adapted-process* ..

sublocale *real-sigma-finite-adapted-process-order* \subseteq *real-sigma-finite-adapted-process* ..

..

locale *sigma-finite-adapted-process-linorder* = *sigma-finite-adapted-process-order*
M F t₀ X for M F t₀ and X :: - ⇒ - ⇒ - :: {linorder-topology}

locale *nat-sigma-finite-adapted-process-linorder* = *sigma-finite-adapted-process-linorder*
M F 0 :: nat X for M F X
locale *real-sigma-finite-adapted-process-linorder* = *sigma-finite-adapted-process-linorder*
M F 0 :: real X for M F X

sublocale *nat-sigma-finite-adapted-process-linorder* ⊆ *nat-sigma-finite-adapted-process-order*
 ..
sublocale *real-sigma-finite-adapted-process-linorder* ⊆ *real-sigma-finite-adapted-process-order*
 ..

9.1 Martingale

locale *martingale* = *sigma-finite-adapted-process* +
assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
and *martingale-property*: $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE \xi \text{ in } M. X i \xi =$
cond-exp M (F i) (X j) ξ

locale *martingale-order* = *martingale M F t₀ X for M F t₀ and X :: - ⇒ - ⇒ -*
:: {order-topology, ordered-real-vector}
locale *martingale-linorder* = *martingale M F t₀ X for M F t₀ and X :: - ⇒ - ⇒ -*
- :: {linorder-topology, ordered-real-vector}
sublocale *martingale-linorder* ⊆ *martingale-order* ..

lemma (**in** *sigma-finite-filtered-measure*) *martingale-const-fun*[*intro*]:
assumes *integrable M f f* ∈ *borel-measurable (F t₀)*
shows *martingale M F t₀ (λ-. f)*
using *assms sigma-finite-subalgebra.cond-exp-F-meas*[*OF - assms(1)*], *THEN AE-symmetric*
borel-measurable-mono
by (*unfold-locale*) *blast+*

lemma (**in** *sigma-finite-filtered-measure*) *martingale-cond-exp*[*intro*]:
assumes *integrable M f*
shows *martingale M F t₀ (λi. cond-exp M (F i) f)*
using *sigma-finite-subalgebra.borel-measurable-cond-exp'* *borel-measurable-cond-exp*

by (*unfold-locale*) (*auto intro: sigma-finite-subalgebra.cond-exp-nested-subalg*[*OF*
- assms] *simp add: subalgebra-F subalgebra*)

corollary (**in** *sigma-finite-filtered-measure*) *martingale-zero*[*intro*]: *martingale M*
F t₀ (λ-. 0) **by** *fastforce*

corollary (**in** *finite-filtered-measure*) *martingale-const*[*intro*]: *martingale M F t₀*
(λ-. c) **by** *fastforce*

9.2 Submartingale

locale *submartingale* = *sigma-finite-adapted-process-order* +
assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M \ (X \ i)$
and *submartingale-property*: $\bigwedge i \ j. t_0 \leq i \implies i \leq j \implies AE \ \xi \text{ in } M. X \ i \ \xi \leq$
cond-exp $M \ (F \ i) \ (X \ j) \ \xi$

locale *submartingale-linorder* = *submartingale* $M \ F \ t_0 \ X$ **for** $M \ F \ t_0$ **and** $X :: -$
 $\Rightarrow - \Rightarrow - :: \{\text{linorder-topology}\}$

sublocale *martingale-order* \subseteq *submartingale* **using** *martingale-property* **by** (*unfold-locales*)
(*force simp add: integrable*) +
sublocale *martingale-linorder* \subseteq *submartingale-linorder* ..

9.3 Supermartingale

locale *supermartingale* = *sigma-finite-adapted-process-order* +
assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M \ (X \ i)$
and *supermartingale-property*: $\bigwedge i \ j. t_0 \leq i \implies i \leq j \implies AE \ \xi \text{ in } M. X \ i \ \xi$
 $\geq \text{cond-exp } M \ (F \ i) \ (X \ j) \ \xi$

locale *supermartingale-linorder* = *supermartingale* $M \ F \ t_0 \ X$ **for** $M \ F \ t_0$ **and** X
 $:: - \Rightarrow - \Rightarrow - :: \{\text{linorder-topology}\}$

sublocale *martingale-order* \subseteq *supermartingale* **using** *martingale-property* **by** (*unfold-locales*)
(*force simp add: integrable*) +
sublocale *martingale-linorder* \subseteq *supermartingale-linorder* ..

lemma *martingale-iff*:

shows *martingale* $M \ F \ t_0 \ X \longleftrightarrow \text{submartingale } M \ F \ t_0 \ X \wedge \text{supermartingale } M \ F \ t_0 \ X$

proof (*rule iffI*)

assume *asm*: *martingale* $M \ F \ t_0 \ X$

interpret *martingale-order* $M \ F \ t_0 \ X$ **by** (*intro martingale-order.intro asm*)

show *submartingale* $M \ F \ t_0 \ X \wedge \text{supermartingale } M \ F \ t_0 \ X$ **using** *submartingale-axioms supermartingale-axioms* **by** *blast*

next

assume *asm*: *submartingale* $M \ F \ t_0 \ X \wedge \text{supermartingale } M \ F \ t_0 \ X$

interpret *submartingale* $M \ F \ t_0 \ X$ **by** (*simp add: asm*)

interpret *supermartingale* $M \ F \ t_0 \ X$ **by** (*simp add: asm*)

show *martingale* $M \ F \ t_0 \ X$ **using** *submartingale-property supermartingale-property*
by (*unfold-locales*) (*intro integrable, blast, force*)

qed

9.4 Martingale Lemmas

context *martingale*

begin

lemma *set-integral-eq*:

assumes $A \in F \ i \ t_0 \leq i \ i \leq j$
shows $\text{set-lebesgue-integral } M \ A \ (X \ i) = \text{set-lebesgue-integral } M \ A \ (X \ j)$
proof –
interpret $\text{sigma-finite-subalgebra } M \ F \ i$ **using** $\text{assms}(2)$ **by** blast
have $\int x \in A. X \ i \ x \ \partial M = \int x \in A. \text{cond-exp } M \ (F \ i) \ (X \ j) \ x \ \partial M$ **using** $\text{martingale-property}[OF \ \text{assms}(2,3)] \ \text{borel-measurable-cond-exp}' \ \text{assms} \ \text{subalgebra} \ \text{subalgebra-def}$ **by** $(\text{intro } \text{set-lebesgue-integral-cong-AE}[OF \ - \ \text{random-variable}]) \ \text{fastforce} +$
also have $\dots = \int x \in A. X \ j \ x \ \partial M$ **using** assms **by** $(\text{auto simp: integrable intro: cond-exp-set-integral[symmetric]})$
finally show $?thesis$.
qed

lemma $\text{scaleR-const}[\text{intro}]$:
shows $\text{martingale } M \ F \ t_0 \ (\lambda i \ x. \ c \ *_R \ X \ i \ x)$
proof –
 $\{$
fix $i \ j :: 'b$ **assume** $\text{asm: } t_0 \leq i \ i \leq j$
interpret $\text{sigma-finite-subalgebra } M \ F \ i$ **using** asm **by** blast
have $\text{AE } x \ \text{in } M. \ c \ *_R \ X \ i \ x = \text{cond-exp } M \ (F \ i) \ (\lambda x. \ c \ *_R \ X \ j \ x) \ x$ **using** $\text{asm} \ \text{cond-exp-scaleR-right}[OF \ \text{integrable, of } j, \ \text{THEN } \text{AE-symmetric}] \ \text{martingale-property}[OF \ \text{asm}]$ **by** force
 $\}$
thus $?thesis$ **by** $(\text{unfold-locales}) \ (\text{auto simp add: integrable martingale.integrable})$
qed

lemma $\text{uminus}[\text{intro}]$:
shows $\text{martingale } M \ F \ t_0 \ (- \ X)$
using $\text{scaleR-const}[of \ -1]$ **by** $(\text{force intro: back-subst}[of \ \text{martingale } M \ F \ t_0])$

lemma $\text{add}[\text{intro}]$:
assumes $\text{martingale } M \ F \ t_0 \ Y$
shows $\text{martingale } M \ F \ t_0 \ (\lambda i \ \xi. \ X \ i \ \xi + Y \ i \ \xi)$
proof –
interpret $Y: \text{martingale } M \ F \ t_0 \ Y$ **by** $(\text{rule } \text{assms})$
 $\{$
fix $i \ j :: 'b$ **assume** $\text{asm: } t_0 \leq i \ i \leq j$
hence $\text{AE } \xi \ \text{in } M. \ X \ i \ \xi + Y \ i \ \xi = \text{cond-exp } M \ (F \ i) \ (\lambda x. \ X \ j \ x + Y \ j \ x) \ \xi$
using $\text{sigma-finite-subalgebra.cond-exp-add}[OF \ - \ \text{integrable martingale.integrable}[OF \ \text{assms}], \ \text{of } F \ i \ j \ j, \ \text{THEN } \text{AE-symmetric}]$
 $\text{martingale-property}[OF \ \text{asm}] \ \text{martingale.martingale-property}[OF \ \text{assms} \ \text{asm}]$ **by** force
 $\}$
thus $?thesis$ **using** assms
by $(\text{unfold-locales}) \ (\text{auto simp add: integrable martingale.integrable})$
qed

lemma $\text{diff}[\text{intro}]$:
assumes $\text{martingale } M \ F \ t_0 \ Y$
shows $\text{martingale } M \ F \ t_0 \ (\lambda i \ x. \ X \ i \ x - Y \ i \ x)$

```

proof –
  interpret  $Y$ : martingale  $M$   $F$   $t_0$   $Y$  by (rule assms)
  {
    fix  $i\ j :: 'b$  assume  $asm$ :  $t_0 \leq i \leq j$ 
    hence  $AE\ \xi$  in  $M$ .  $X\ i\ \xi - Y\ i\ \xi = \text{cond-exp } M\ (F\ i)\ (\lambda x. X\ j\ x - Y\ j\ x)\ \xi$ 
    using sigma-finite-subalgebra.cond-exp-diff[OF - integrable martingale.integrable[OF
assms], of  $F\ i\ j$ , THEN AE-symmetric]
    martingale-property[OF asm] martingale.martingale-property[OF assms
asm] by fastforce
  }
  thus ?thesis using assms by (unfold-locales) (auto simp add: integrable martin-
gale.integrable)
qed

end

lemma (in sigma-finite-adapted-process) martingale-of-set-integral-eq:
  assumes integrable:  $\bigwedge i. \text{integrable } M\ (X\ i)$ 
  and  $\bigwedge A\ i\ j. t_0 \leq i \implies i \leq j \implies A \in F\ i \implies \text{set-lebesgue-integral } M\ A\ (X\ i) = \text{set-lebesgue-integral } M\ A\ (X\ j)$ 
  shows martingale  $M$   $F$   $t_0$   $X$ 
proof (unfold-locales)
  fix  $i\ j :: 'b$  assume  $asm$ :  $t_0 \leq i \leq j$ 
  interpret sigma-finite-subalgebra  $M$   $F$   $i$  using  $asm$  by blast
  interpret  $r$ : sigma-finite-measure restr-to-subalg  $M$  ( $F\ i$ ) by (simp add: sigma-fin-subalg)
  {
    fix  $A$  assume  $A \in \text{restr-to-subalg } M\ (F\ i)$ 
    hence  $*$ :  $A \in F\ i$  using sets-restr-to-subalg subalgebra  $asm$  by blast
    have set-lebesgue-integral (restr-to-subalg  $M$  ( $F\ i$ ))  $A$  ( $X\ i$ ) = set-lebesgue-integral
 $M\ A$  ( $X\ i$ ) using  $*$  subalg  $asm$  by (auto simp: set-lebesgue-integral-def intro: inte-
gral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator)
    also have  $\dots = \text{set-lebesgue-integral } M\ A\ (\text{cond-exp } M\ (F\ i)\ (X\ j))$  using  $*$ 
assms(2)[OF  $asm$ ] cond-exp-set-integral[OF integrable] by auto
    finally have set-lebesgue-integral (restr-to-subalg  $M$  ( $F\ i$ ))  $A$  ( $X\ i$ ) = set-lebesgue-integral
(restr-to-subalg  $M$  ( $F\ i$ ))  $A$  (cond-exp  $M$  ( $F\ i$ ) ( $X\ j$ )) using  $*$  subalg by (auto simp:
set-lebesgue-integral-def intro!: integral-subalgebra2[symmetric] borel-measurable-scaleR
borel-measurable-cond-exp borel-measurable-indicator)
  }
  hence  $AE\ \xi$  in restr-to-subalg  $M$  ( $F\ i$ ).  $X\ i\ \xi = \text{cond-exp } M\ (F\ i)\ (X\ j)\ \xi$ 
using  $asm$  by (intro r.density-unique, auto intro: integrable-in-subalg subalg
borel-measurable-cond-exp integrable)
  thus  $AE\ \xi$  in  $M$ .  $X\ i\ \xi = \text{cond-exp } M\ (F\ i)\ (X\ j)\ \xi$  using AE-restr-to-subalg[OF
subalg] by blast
qed (simp add: integrable)

```

9.5 Submartingale Lemmas

```

context submartingale
begin

```

lemma *cond-exp-diff-nonneg*:
assumes $t_0 \leq i \leq j$
shows $AE\ x\ in\ M. 0 \leq cond_exp\ M\ (F\ i)\ (\lambda \xi. X\ j\ \xi - X\ i\ \xi)\ x$
using *submartingale-property*[*OF assms*] *assms sigma-finite-subalgebra.cond-exp-diff*[*OF*
- *integrable*(*1,1*), *of - j i*] *sigma-finite-subalgebra.cond-exp-F-meas*[*OF - integrable*
adapted, of i] **by** *fastforce*

lemma *add*[*intro*]:
assumes *submartingale* $M\ F\ t_0\ Y$
shows *submartingale* $M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi + Y\ i\ \xi)$
proof -
interpret $Y: submartingale\ M\ F\ t_0\ Y$ **by** (*rule assms*)
{
fix $i\ j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
hence $AE\ \xi\ in\ M. X\ i\ \xi + Y\ i\ \xi \leq cond_exp\ M\ (F\ i)\ (\lambda x. X\ j\ x + Y\ j\ x)\ \xi$
using *sigma-finite-subalgebra.cond-exp-add*[*OF - integrable submartingale.integrable*[*OF*
assms], *of F i j j*]
submartingale-property[*OF asm*] *submartingale.submartingale-property*[*OF*
assms asm] *add-mono*[*of X i - - Y i -*] **by** *force*
}
thus *?thesis using assms by* (*unfold-locals*) (*auto simp add: borel-measurable-add*
random-variable adapted integrable Y.random-variable Y.adapted submartingale.integrable)

qed

lemma *diff*[*intro*]:
assumes *supermartingale* $M\ F\ t_0\ Y$
shows *submartingale* $M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi - Y\ i\ \xi)$
proof -
interpret $Y: supermartingale\ M\ F\ t_0\ Y$ **by** (*rule assms*)
{
fix $i\ j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
hence $AE\ \xi\ in\ M. X\ i\ \xi - Y\ i\ \xi \leq cond_exp\ M\ (F\ i)\ (\lambda x. X\ j\ x - Y\ j\ x)\ \xi$
using *sigma-finite-subalgebra.cond-exp-diff*[*OF - integrable supermartingale.integrable*[*OF*
assms], *of F i j j*]
submartingale-property[*OF asm*] *supermartingale.supermartingale-property*[*OF*
assms asm] *diff-mono*[*of X i - - Y i -*] **by** *force*
}
thus *?thesis using assms by* (*unfold-locals*) (*auto simp add: borel-measurable-diff*
random-variable adapted integrable Y.random-variable Y.adapted supermartingale.integrable)

qed

lemma *scaleR-nonneg*:
assumes $c \geq 0$
shows *submartingale* $M\ F\ t_0\ (\lambda i\ \xi. c *_{\mathbb{R}} X\ i\ \xi)$
proof
{


```

    fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
    thus AE  $\xi$  in  $M$ .  $c *_R X i \xi \leq \text{cond-exp } M (F i) (\lambda \xi. c *_R X j \xi) \xi$ 
    using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of F i
j c] submartingale-property[OF asm] by (fastforce intro!: scaleR-left-mono[OF -
assms])
  }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)

```

lemma *scaleR-nonpos*:

```

  assumes  $c \leq 0$ 
  shows supermartingale  $M F t_0 (\lambda i \xi. c *_R X i \xi)$ 
proof
  {
    fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
    thus AE  $\xi$  in  $M$ .  $c *_R X i \xi \geq \text{cond-exp } M (F i) (\lambda \xi. c *_R X j \xi) \xi$ 
    using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of F i j
c] submartingale-property[OF asm]
    by (fastforce intro!: scaleR-left-mono-neg[OF - assms])
  }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)

```

lemma *uminus[intro]*:

```

  shows supermartingale  $M F t_0 (- X)$ 
  unfolding fun-Comp-def using scaleR-nonpos[of -1] by simp

```

end

context *submartingale-linorder*
begin

lemma *set-integral-le*:

```

  assumes  $A \in F i t_0 \leq i \leq j$ 
  shows set-lebesgue-integral  $M A (X i) \leq \text{set-lebesgue-integral } M A (X j)$ 
  using submartingale-property[OF assms(2), of j] assms subalgebra
  by (subst sigma-finite-subalgebra.cond-exp-set-integral[OF - integrable assms(1),
of j])
  (auto intro!: scaleR-left-mono integral-mono-AE-banach integrable-mult-indicator
integrable simp add: subalgebra-def set-lebesgue-integral-def)

```

lemma *max*:

```

  assumes submartingale-linorder  $M F t_0 Y$ 
  shows submartingale-linorder  $M F t_0 (\lambda i \xi. \max (X i \xi) (Y i \xi))$ 
proof (unfold-locales)
  interpret  $Y$ : submartingale-linorder  $M F t_0 Y$  by (rule assms)
  {
    fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
    have AE  $\xi$  in  $M$ .  $\max (X i \xi) (Y i \xi) \leq \max (\text{cond-exp } M (F i) (X j) \xi)$ 

```

(cond-exp $M (F i) (Y j) \xi$) **using** *submartingale-property* Y .*submartingale-property*
asm **unfolding** *max-def* **by** *fastforce*
 thus $AE \xi$ in M . $\max (X i \xi) (Y i \xi) \leq \text{cond-exp } M (F i) (\lambda \xi. \max (X j \xi) (Y j \xi)) \xi$ **using** *sigma-finite-subalgebra.cond-exp-max*[*OF* - *integrable* Y .*integrable*, of $F i j j$] *asm* **by** (*fast intro*: *order.trans*)
 }
 show $\bigwedge i. t_0 \leq i \implies (\lambda \xi. \max (X i \xi) (Y i \xi)) \in \text{borel-measurable } (F i) \bigwedge i. t_0 \leq i \implies \text{integrable } M (\lambda \xi. \max (X i \xi) (Y i \xi))$ **by** (*force intro*: Y .*integrable integrable assms*)
qed

lemma *max-0*:

shows *submartingale-linorder* $M F t_0 (\lambda i \xi. \max 0 (X i \xi))$
proof –
 interpret *zero*: *martingale-linorder* $M F t_0 \lambda - . 0$ **by** (*force intro*: *martingale-linorder.intro* *martingale-order.intro*)
 show ?thesis **by** (*intro zero.max submartingale-linorder.intro* *submartingale-axioms*)
qed
end

lemma (in *sigma-finite-adapted-process-order*) *submartingale-of-cond-exp-diff-nonneg*:

assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
 and *diff-nonneg*: $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE x$ in $M. 0 \leq \text{cond-exp } M (F i) (\lambda \xi. X j \xi - X i \xi) x$
 shows *submartingale* $M F t_0 X$
proof (*unfold-locales*)
 {
 fix $i j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
 thus $AE \xi$ in $M. X i \xi \leq \text{cond-exp } M (F i) (X j) \xi$
using *diff-nonneg*[*OF* *asm*] *sigma-finite-subalgebra.cond-exp-diff*[*OF* - *integrable*(1,1), of $F i j i$]
sigma-finite-subalgebra.cond-exp-F-meas[*OF* - *integrable adapted*, of i] **by** *fastforce*
 }
qed (*intro integrable*)

lemma (in *sigma-finite-adapted-process-linorder*) *submartingale-of-set-integral-le*:

assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
 and $\bigwedge A i j. t_0 \leq i \implies i \leq j \implies A \in F i \implies \text{set-lebesgue-integral } M A (X i) \leq \text{set-lebesgue-integral } M A (X j)$
 shows *submartingale* $M F t_0 X$
proof (*unfold-locales*)
 {
 fix $i j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
 interpret r : *sigma-finite-measure restr-to-subalg* $M (F i)$ **using** *asm sigma-finite-subalgebra.sigma-fin-subalg*
by *blast*
 {
 fix A **assume** $A \in \text{restr-to-subalg } M (F i)$

hence *: $A \in F i$ **using** *asm sets-restr-to-subalg subalgebra* **by** *blast*
have *set-lebesgue-integral (restr-to-subalg M (F i)) A (X i) = set-lebesgue-integral M A (X i)* **using** * *asm subalgebra* **by** (*auto simp: set-lebesgue-integral-def intro: integral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator*)
also have $\dots \leq \text{set-lebesgue-integral } M A (\text{cond-exp } M (F i) (X j))$ **using**
* *assms(2)[OF asm] asm sigma-finite-subalgebra.cond-exp-set-integral[OF - integrable]* **by** *fastforce*
also have $\dots = \text{set-lebesgue-integral (restr-to-subalg } M (F i)) A (\text{cond-exp } M (F i) (X j))$ **using** * *asm subalgebra* **by** (*auto simp: set-lebesgue-integral-def intro!: integral-subalgebra2[symmetric] borel-measurable-scaleR borel-measurable-cond-exp borel-measurable-indicator*)
finally have $0 \leq \text{set-lebesgue-integral (restr-to-subalg } M (F i)) A (\lambda \xi. \text{cond-exp } M (F i) (X j) \xi - X i \xi)$ **using** * *asm subalgebra* **by** (*subst set-integral-diff, auto simp add: set-integrable-def sets-restr-to-subalg intro!: integrable adapted integrable-in-subalg borel-measurable-scaleR borel-measurable-indicator borel-measurable-cond-exp integrable-mult-indicator*)
}
hence $AE \xi \text{ in } \text{restr-to-subalg } M (F i). 0 \leq \text{cond-exp } M (F i) (X j) \xi - X i \xi$
by (*intro r.density-nonneg integrable-in-subalg asm subalgebra borel-measurable-diff borel-measurable-cond-exp adapted Bochner-Integration.integrable-diff integrable-cond-exp integrable*)
thus $AE \xi \text{ in } M. X i \xi \leq \text{cond-exp } M (F i) (X j) \xi$ **using** *AE-restr-to-subalg[OF subalgebra] asm* **by** *simp*
}
qed (*intro integrable*)

9.6 Supermartingale Lemmas

The following lemmas are exact duals of the ones for submartingales.

context *supermartingale*
begin

lemma *cond-exp-diff-nonneg*:
assumes $t_0 \leq i \leq j$
shows $AE x \text{ in } M. 0 \leq \text{cond-exp } M (F i) (\lambda \xi. X i \xi - X j \xi) x$
using *assms supermartingale-property[OF assms] sigma-finite-subalgebra.cond-exp-diff[OF - integrable(1,1), of F i i j]*
sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] **by**
fastforce

lemma *add[intro]*:
assumes *supermartingale M F t_0 Y*
shows *supermartingale M F t_0 ($\lambda i \xi. X i \xi + Y i \xi$)*
proof –
interpret $Y: \text{supermartingale } M F t_0 Y$ **by** (*rule assms*)
{
fix $i j :: 'b$ **assume** $asm: t_0 \leq i \leq j$
hence $AE \xi \text{ in } M. X i \xi + Y i \xi \geq \text{cond-exp } M (F i) (\lambda x. X j x + Y j x) \xi$
using *sigma-finite-subalgebra.cond-exp-add[OF - integrable supermartingale.integrable[OF*

```

assms], of F i j j]
  supermartingale-property[OF asm] supermartingale.supermartingale-property[OF
assms asm] add-mono[of - X i - - Y i -] by force
}
thus ?thesis using assms by (unfold-locale) (auto simp add: borel-measurable-add
random-variable adapted integrable Y.random-variable Y.adapted supermartingale.integrable)

qed

```

```

lemma diff[intro]:
  assumes submartingale M F t0 Y
  shows supermartingale M F t0 (λi ξ. X i ξ - Y i ξ)
proof -
  interpret Y: submartingale M F t0 Y by (rule assms)
  {
    fix i j :: 'b assume asm: t0 ≤ i i ≤ j
    hence AE ξ in M. X i ξ - Y i ξ ≥ cond-exp M (F i) (λx. X j x - Y j x) ξ
    using sigma-finite-subalgebra.cond-exp-diff[OF - integrable submartingale.integrable[OF
assms], of F i j j, unfolded fun-diff-def]
    supermartingale-property[OF asm] submartingale.submartingale-property[OF
assms asm] diff-mono[of - X i - Y i -] by force
  }
  thus ?thesis using assms by (unfold-locale) (auto simp add: borel-measurable-diff
random-variable adapted integrable Y.random-variable Y.adapted submartingale.integrable)

```

qed

```

lemma scaleR-nonneg:
  assumes c ≥ 0
  shows supermartingale M F t0 (λi ξ. c *R X i ξ)
proof
  {
    fix i j :: 'b assume asm: t0 ≤ i i ≤ j
    thus AE ξ in M. c *R X i ξ ≥ cond-exp M (F i) (λξ. c *R X j ξ) ξ
    using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of F i
j c] supermartingale-property[OF asm] by (fastforce intro!: scaleR-left-mono[OF -
assms])
  }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)

```

```

lemma scaleR-nonpos:
  assumes c ≤ 0
  shows submartingale M F t0 (λi ξ. c *R X i ξ)
proof
  {
    fix i j :: 'b assume asm: t0 ≤ i i ≤ j
    thus AE ξ in M. c *R X i ξ ≤ cond-exp M (F i) (λξ. c *R X j ξ) ξ
    using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of F i j c]

```

supermartingale-property[*OF asm*] **by** (*fastforce intro!*: *scaleR-left-mono-neg*[*OF - assms*])
 }
qed (*auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable random-variable adapted borel-measurable-const-scaleR*)

lemma *uminus*[*intro*]:
 shows *submartingale* *M F t₀* ($- X$)
 unfolding *fun-Compl-def* **using** *scaleR-nonpos*[*of -1*] **by** *simp*

end

context *supermartingale-linorder*
begin

lemma *set-integral-ge*:
 assumes $A \in F$ $i t_0 \leq i i \leq j$
 shows *set-lebesgue-integral* *M A* ($X i$) \geq *set-lebesgue-integral* *M A* ($X j$)
 using *supermartingale-property*[*OF assms*(2), *of j*] *assms subalgebra*
 by (*subst sigma-finite-subalgebra.cond-exp-set-integral*[*OF - integrable assms*(1), *of j*])
 (*auto intro!*: *scaleR-left-mono integral-mono-AE-banach integrable-mult-indicator integrable simp add: subalgebra-def set-lebesgue-integral-def*)

lemma *min*:
 assumes *supermartingale-linorder* *M F t₀* *Y*
 shows *supermartingale-linorder* *M F t₀* ($\lambda i \xi. \min (X i \xi) (Y i \xi)$)
proof (*unfold-locales*)
 interpret *Y*: *supermartingale-linorder* *M F t₀* *Y* **by** (*rule assms*)
 {
 fix $i j :: 'b$ **assume** *asm*: $t_0 \leq i i \leq j$
 have $AE \xi \text{ in } M. \min (X i \xi) (Y i \xi) \geq \min (\text{cond-exp } M (F i) (X j) \xi) (\text{cond-exp } M (F i) (Y j) \xi)$ **using** *supermartingale-property Y.supermartingale-property asm*
 unfolding *min-def* **by** *fastforce*
 thus $AE \xi \text{ in } M. \min (X i \xi) (Y i \xi) \geq \text{cond-exp } M (F i) (\lambda \xi. \min (X j \xi) (Y j \xi)) \xi$ **using** *sigma-finite-subalgebra.cond-exp-min*[*OF - integrable Y.integrable, of F i j j*] *asm* **by** (*fast intro: order.trans*)
 }
 show $\bigwedge i. t_0 \leq i \implies (\lambda \xi. \min (X i \xi) (Y i \xi)) \in \text{borel-measurable } (F i) \bigwedge i. t_0 \leq i \implies \text{integrable } M (\lambda \xi. \min (X i \xi) (Y i \xi))$ **by** (*force intro: Y.integrable integrable assms*)
qed

lemma *min-0*:
 shows *supermartingale-linorder* *M F t₀* ($\lambda i \xi. \min 0 (X i \xi)$)
proof –
 interpret *zero*: *martingale-linorder* *M F t₀* $\lambda - . 0$ **by** (*force intro: martingale-linorder.intro*)
 show ?thesis **by** (*intro zero.min supermartingale-linorder.intro supermartin-*

gale-axioms)
qed

end

lemma (in *sigma-finite-adapted-process-order*) *supermartingale-of-cond-exp-diff-nonneg*:

assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M \ (X \ i)$
and *diff-nonneg*: $\bigwedge i \ j. t_0 \leq i \implies i \leq j \implies AE \ x \text{ in } M. \ 0 \leq \text{cond-exp } M \ (F \ i) \ (\lambda \xi. X \ i \ \xi - X \ j \ \xi) \ x$
shows *supermartingale* $M \ F \ t_0 \ X$
proof
 $\{$
fix $i \ j :: 'b$ **assume** $asm: t_0 \leq i \ i \leq j$
thus $AE \ \xi \text{ in } M. X \ i \ \xi \geq \text{cond-exp } M \ (F \ i) \ (X \ j) \ \xi$
using *diff-nonneg*[*OF asm*] *sigma-finite-subalgebra.cond-exp-diff*[*OF - integrable(1,1), of F i i j*]
sigma-finite-subalgebra.cond-exp-F-meas[*OF - integrable adapted, of i*] **by**
fastforce
 $\}$
qed (*intro integrable*)

lemma (in *sigma-finite-adapted-process-linorder*) *supermartingale-of-set-integral-ge*:

assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M \ (X \ i)$
and $\bigwedge A \ i \ j. t_0 \leq i \implies i \leq j \implies A \in F \ i \implies \text{set-lebesgue-integral } M \ A \ (X \ i) \leq \text{set-lebesgue-integral } M \ A \ (X \ j)$
shows *supermartingale* $M \ F \ t_0 \ X$
proof –
interpret $-:$ *adapted-process* $M \ F \ t_0 \ -X$ **by** (*rule uminus*)
interpret *uminus-X*: *sigma-finite-adapted-process-linorder* $M \ F \ t_0 \ -X$..
note $*$ = *set-integral-uminus*[*unfolded set-integrable-def, OF integrable-mult-indicator*[*OF - integrable*]]
have *supermartingale* $M \ F \ t_0 \ (-(-X))$
using *ord-eq-le-trans*[*OF * ord-le-eq-trans*[*OF le-imp-neg-le*[*OF assms(2)*] * [*symmetric*]]]
subalgebra
by (*intro submartingale.uminus uminus-X.submartingale-of-set-integral-le*)
(*clarsimp simp add: fun-Compl-def subalgebra-def integrable | fastforce*) +
thus *?thesis* **unfolding** *fun-Compl-def* **by** *simp*
qed

9.7 Discrete Time Martingales

locale *nat-martingale* = *martingale* $M \ F \ 0 :: \text{nat } X$ **for** $M \ F \ X$

locale *nat-submartingale* = *submartingale* $M \ F \ 0 :: \text{nat } X$ **for** $M \ F \ X$

locale *nat-supermartingale* = *supermartingale* $M \ F \ 0 :: \text{nat } X$ **for** $M \ F \ X$

locale *nat-submartingale-linorder* = *submartingale-linorder* $M \ F \ 0 :: \text{nat } X$ **for** $M \ F \ X$

locale *nat-supermartingale-linorder* = *supermartingale-linorder* $M \ F \ 0 :: \text{nat } X$

for $M F X$

sublocale $\text{nat-submartingale-linorder} \subseteq \text{nat-submartingale} \dots$

sublocale $\text{nat-supermartingale-linorder} \subseteq \text{nat-supermartingale} \dots$

9.8 Discrete Time Martingales

lemma (in nat-martingale) *predictable-eq-zero*:

assumes $\text{nat-predictable-process } M F X$

shows $\text{AE } \xi \text{ in } M. X \ i \ \xi = X \ 0 \ \xi$

proof (induction i)

case 0

then show $?case$ **by** (simp add: bot-nat-def)

next

case (Suc i)

interpret S : $\text{nat-adapted-process } M F \lambda i. X \ (Suc \ i)$ **by** (intro $\text{nat-predictable-process.adapted-Suc assms}$)

show $?case$ **using** $Suc \ S.adapted[of \ i] \ \text{martingale-property}[OF \ - \ le-SucI, \ of \ i] \ \text{sigma-finite-subalgebra.cond-exp-F-meas}[OF \ - \ integrable, \ of \ F \ i \ Suc \ i]$ **by** fastforce

qed

lemma (in $\text{nat-sigma-finite-adapted-process}$) *martingale-of-set-integral-eq-Suc*:

assumes $\text{integrable: } \bigwedge i. \text{integrable } M \ (X \ i)$

and $\bigwedge A \ i. A \in F \ i \implies \text{set-lebesgue-integral } M \ A \ (X \ i) = \text{set-lebesgue-integral } M \ A \ (X \ (Suc \ i))$

shows $\text{nat-martingale } M F X$

proof (intro $\text{nat-martingale.intro martingale-of-set-integral-eq}$)

fix $i \ j \ A$ **assume** $asm: i \leq j \ A \in \text{sets } (F \ i)$

show $\text{set-lebesgue-integral } M \ A \ (X \ i) = \text{set-lebesgue-integral } M \ A \ (X \ j)$ **using** asm

proof (induction $j - i$ arbitrary: $i \ j$)

case 0

then show $?case$ **using** asm **by** simp

next

case (Suc n)

hence $*$: $n = j - Suc \ i$ **by** linarith

have $Suc \ i \leq j$ **using** $Suc(2,3)$ **by** linarith

thus $?case$ **using** $\text{sets-F-mono}[OF \ - \ le-SucI] \ Suc(4) \ Suc(1)[OF \ *]$ **by** (auto intro: $assms(2)[THEN \ trans]$)

qed

qed (simp add: integrable)

lemma (in $\text{nat-sigma-finite-adapted-process}$) *martingale-nat*:

assumes $\text{integrable: } \bigwedge i. \text{integrable } M \ (X \ i)$

and $\bigwedge i. \text{AE } \xi \text{ in } M. X \ i \ \xi = \text{cond-exp } M \ (F \ i) \ (X \ (Suc \ i)) \ \xi$

shows $\text{nat-martingale } M F X$

proof (unfold-locales)

fix $i \ j :: \text{nat}$ **assume** $asm: i \leq j$

show $\text{AE } \xi \text{ in } M. X \ i \ \xi = \text{cond-exp } M \ (F \ i) \ (X \ j) \ \xi$ **using** asm

```

proof (induction  $j - i$  arbitrary:  $i\ j$ )
  case 0
    hence  $j = i$  by simp
    thus ?case using sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted,
    THEN AE-symmetric] by blast
  next
    case (Suc  $n$ )
    have  $j: j = \text{Suc } (n + i)$  using Suc by linarith
    have  $n: n = n + i - i$  using Suc by linarith
    have *: AE  $\xi$  in  $M$ . cond-exp  $M$  ( $F$  ( $n + i$ )) ( $X\ j$ )  $\xi = X$  ( $n + i$ )  $\xi$  unfolding
     $j$  using assms(2)[THEN AE-symmetric] by blast
    have AE  $\xi$  in  $M$ . cond-exp  $M$  ( $F\ i$ ) ( $X\ j$ )  $\xi = \text{cond-exp } M$  ( $F\ i$ ) ( $\text{cond-exp } M$ 
    ( $F$  ( $n + i$ )) ( $X\ j$ ))  $\xi$  by (intro cond-exp-nested-subalg integrable subalg, simp add:
    subalgebra-def space-F sets-F-mono)
    hence AE  $\xi$  in  $M$ . cond-exp  $M$  ( $F\ i$ ) ( $X\ j$ )  $\xi = \text{cond-exp } M$  ( $F\ i$ ) ( $X$  ( $n + i$ ))
     $\xi$  using cond-exp-cong-AE[OF integrable-cond-exp integrable *] by force
    thus ?case using Suc(1)[OF  $n$ ] by fastforce
  qed
qed (simp add: integrable)

```

```

lemma (in nat-sigma-finite-adapted-process) martingale-of-cond-exp-diff-Suc-eq-zero:
  assumes integrable:  $\bigwedge i$ . integrable  $M$  ( $X\ i$ )
    and  $\bigwedge i$ . AE  $\xi$  in  $M$ .  $0 = \text{cond-exp } M$  ( $F\ i$ ) ( $\lambda \xi. X$  ( $\text{Suc } i$ )  $\xi - X\ i\ \xi$ )  $\xi$ 
    shows nat-martingale  $M\ F\ X$ 
proof (intro martingale-nat integrable)
  fix  $i$ 
  show AE  $\xi$  in  $M$ .  $X\ i\ \xi = \text{cond-exp } M$  ( $F\ i$ ) ( $X$  ( $\text{Suc } i$ ))  $\xi$  using cond-exp-diff[OF
  integrable(1,1), of  $i$  Suc  $i\ i$ ] cond-exp-F-meas[OF integrable adapted, of  $i$ ] assms(2)[of
   $i$ ] by fastforce
qed

```

9.9 Discrete Time Submartingales

```

lemma (in nat-submartingale) predictable-ge-zero:
  assumes nat-predictable-process  $M\ F\ X$ 
  shows AE  $\xi$  in  $M$ .  $X\ i\ \xi \geq X\ 0\ \xi$ 
proof (induction  $i$ )
  case 0
    then show ?case by (simp add: bot-nat-def)
  next
    case (Suc  $i$ )
    interpret  $S$ : nat-adapted-process  $M\ F\ \lambda i. X$  ( $\text{Suc } i$ ) by (intro nat-predictable-process.adapted-Suc
    assms)
    show ?case using Suc  $S$ .adapted[of  $i$ ] submartingale-property[OF - le-SucI, of  $i$ ]
    sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable, of  $F\ i$  Suc  $i$ ] by fastforce
  qed

```

```

lemma (in nat-sigma-finite-adapted-process-linorder) submartingale-of-set-integral-le-Suc:
  assumes integrable:  $\bigwedge i$ . integrable  $M$  ( $X\ i$ )

```


and $\bigwedge i. A \in F i \implies \text{set-lebesgue-integral } M A (X i) \leq \text{set-lebesgue-integral } M A (X (\text{Suc } i))$
shows *nat-submartingale* $M F X$
proof (*intro nat-submartingale.intro submartingale-of-set-integral-le*)
fix $i j A$ **assume** $\text{asm}: i \leq j \ A \in \text{sets } (F i)$
show $\text{set-lebesgue-integral } M A (X i) \leq \text{set-lebesgue-integral } M A (X j)$ **using** *asm*
proof (*induction j - i arbitrary: i j*)
case 0
then show ?*case* **using** *asm* **by** *simp*
next
case (*Suc n*)
hence $*$: $n = j - \text{Suc } i$ **by** *linarith*
have $\text{Suc } i \leq j$ **using** *Suc(2,3)* **by** *linarith*
thus ?*case* **using** *sets-F-mono[OF - le-SucI] Suc(4) Suc(1)[OF *]* **by** (*auto intro: assms(2)[THEN order-trans]*)
qed
qed (*simp add: integrable*)

lemma (*in nat-sigma-finite-adapted-process-linorder*) *submartingale-nat*:
assumes *integrable*: $\bigwedge i. \text{integrable } M (X i)$
and $\bigwedge i. AE \ \xi \text{ in } M. X i \ \xi \leq \text{cond-exp } M (F i) (X (\text{Suc } i)) \ \xi$
shows *nat-submartingale* $M F X$
using *subalg integrable assms(2)*
by (*intro submartingale-of-set-integral-le-Suc ord-le-eq-trans[OF set-integral-mono-AE-banach cond-exp-set-integral[symmetric]], simp*)
(meson in-mono integrable-mult-indicator set-integrable-def subalgebra-def, meson integrable-cond-exp in-mono integrable-mult-indicator set-integrable-def subalgebra-def, fast+)

lemma (*in nat-sigma-finite-adapted-process-linorder*) *submartingale-of-cond-exp-diff-Suc-nonneg*:
assumes *integrable*: $\bigwedge i. \text{integrable } M (X i)$
and $\bigwedge i. AE \ \xi \text{ in } M. 0 \leq \text{cond-exp } M (F i) (\lambda \xi. X (\text{Suc } i) \ \xi - X i \ \xi)$
shows *nat-submartingale* $M F X$
proof (*intro submartingale-nat integrable*)
fix i
show $AE \ \xi \text{ in } M. X i \ \xi \leq \text{cond-exp } M (F i) (X (\text{Suc } i)) \ \xi$ **using** *cond-exp-diff[OF integrable(1,1), of i Suc i i] cond-exp-F-meas[OF integrable adapted, of i] assms(2)[of i]* **by** *fastforce*
qed

lemma (*in nat-submartingale-linorder*) *partial-sum-scaleR*:
assumes *nat-adapted-process* $M F C \ \bigwedge i. AE \ \xi \text{ in } M. 0 \leq C i \ \xi \ \bigwedge i. AE \ \xi \text{ in } M. C i \ \xi \leq R$
shows *nat-submartingale* $M F (\lambda n \ \xi. \sum_{i < n}. C i \ \xi *_{\mathbb{R}} (X (\text{Suc } i) \ \xi - X i \ \xi))$
proof –
interpret C : *nat-adapted-process* $M F C$ **by** (*rule assms*)
interpret C' : *nat-adapted-process* $M F \lambda i \ \xi. C (i - 1) \ \xi *_{\mathbb{R}} (X i \ \xi - X (i - 1) \ \xi)$
by (*intro nat-adapted-process.intro adapted-process.scaleR-right adapted-process.diff*,

unfold-locales) (*auto intro: adaptedD C.adaptedD*) +
interpret C'' : *nat-adapted-process* $M F \lambda n \xi. \sum i < n. C i \xi *_{\mathbb{R}} (X (Suc i) \xi - X i \xi)$ **by** (*rule C'.partial-sum-Suc[unfolded diff-Suc-1]*)
interpret S : *nat-sigma-finite-adapted-process-linorder* $M F (\lambda n \xi. \sum i < n. C i \xi *_{\mathbb{R}} (X (Suc i) \xi - X i \xi))$..
have *integrable* $M (\lambda x. C i x *_{\mathbb{R}} (X (Suc i) x - X i x))$ **for** i **using** *assms(2,3)[of i]* **by** (*intro Bochner-Integration.integrable-bound[OF integrable-scaleR-right, OF Bochner-Integration.integrable-diff, OF integrable(1,1), of R Suc i i]*) (*auto simp add: mult-mono*)
moreover have $AE \xi$ *in* $M. 0 \leq cond\text{-}exp M (F i) (\lambda \xi. (\sum i < Suc i. C i \xi *_{\mathbb{R}} (X (Suc i) \xi - X i \xi)) - (\sum i < i. C i \xi *_{\mathbb{R}} (X (Suc i) \xi - X i \xi))) \xi$ **for** i
using *sigma-finite-subalgebra.cond-exp-measurable-scaleR[OF - calculation - C.adapted, of i]*
cond-exp-diff-nonneg[OF - le-SucI, OF - order.refl, of i] assms(2,3)[of i]
by (*fastforce simp add: scaleR-nonneg-nonneg integrable*)
ultimately show *?thesis* **by** (*intro S.submartingale-of-cond-exp-diff-Suc-nonneg Bochner-Integration.integrable-sum, blast+*)
qed

lemma (*in nat-submartingale-linorder*) *partial-sum-scaleR'*:
assumes *nat-predictable-process* $M F C \wedge i. AE \xi$ *in* $M. 0 \leq C i \xi \wedge i. AE \xi$ *in* $M. C i \xi \leq R$
shows *nat-submartingale* $M F (\lambda n \xi. \sum i < n. C (Suc i) \xi *_{\mathbb{R}} (X (Suc i) \xi - X i \xi))$
proof –
interpret C : *nat-predictable-process* $M F C$ **by** (*rule assms*)
interpret $Suc\text{-}C$: *nat-adapted-process* $M F \lambda i. C (Suc i)$ **using** $C.adapted\text{-}Suc$.
show *?thesis* **by** (*intro partial-sum-scaleR[of - R] assms*) (*intro-locales*)
qed

9.10 Discrete Time Supermartingales

lemma (*in nat-supermartingale*) *predictable-le-zero*:
assumes *nat-predictable-process* $M F X$
shows $AE \xi$ *in* $M. X i \xi \leq X 0 \xi$
proof (*induction i*)
case 0
then show *?case* **by** (*simp add: bot-nat-def*)
next
case ($Suc i$)
interpret S : *nat-adapted-process* $M F \lambda i. X (Suc i)$ **by** (*intro nat-predictable-process.adapted-Suc assms*)
show *?case* **using** $S.adapted[of i]$ *supermartingale-property[OF - le-SucI, of i]*
sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable, of F i Suc i] **by** *fastforce*
qed

lemma (*in nat-sigma-finite-adapted-process-linorder*) *supermartingale-of-set-integral-ge-Suc*:
assumes *integrable*: $\wedge i. integrable M (X i)$
and $\wedge A i. A \in F i \implies set\text{-lebesgue-integral } M A (X (Suc i)) \leq set\text{-lebesgue-integral}$

$M A (X i)$
shows *nat-supermartingale* $M F X$
proof –
interpret \vdash *adapted-process* $M F 0 -X$ **by** (*rule* *uminus*)
interpret *uminus-X*: *nat-sigma-finite-adapted-process-linorder* $M F -X$..
note $\ast = \text{set-integral-uminus}[\text{unfolded set-integrable-def}, OF \text{ integrable-mult-indicator}[OF$
 $- \text{integrable}]]$
have *nat-supermartingale* $M F (-(-X))$
using *ord-eq-le-trans* $[OF \ast \text{ord-le-eq-trans}[OF \text{le-imp-neg-le}[OF \text{assms}(2)] \ast [\text{symmetric}]]]$
subalgebra
by (*intro* *nat-supermartingale.intro* *submartingale.uminus* *nat-submartingale.axioms*
uminus-X.submartingale-of-set-integral-le-Suc)
(clarsimp simp add: fun-Compl-def subalgebra-def integrable | fastforce) +
thus *?thesis* **unfolding** *fun-Compl-def* **by** *simp*
qed

lemma (*in* *nat-sigma-finite-adapted-process-linorder*) *supermartingale-nat*:
assumes *integrable*: $\bigwedge i. \text{integrable } M (X i)$
and $\bigwedge i. AE \xi \text{ in } M. X i \xi \geq \text{cond-exp } M (F i) (X (Suc i)) \xi$
shows *nat-supermartingale* $M F X$
proof –
interpret \vdash *adapted-process* $M F 0 -X$ **by** (*rule* *uminus*)
interpret *uminus-X*: *nat-sigma-finite-adapted-process-linorder* $M F -X$..
have $AE \xi \text{ in } M. -X i \xi \leq \text{cond-exp } M (F i) (\lambda x. -X (Suc i) x) \xi$ **for** i **using**
assms(2) cond-exp-uminus[OF integrable, of i Suc i] **by** *force*
hence *nat-supermartingale* $M F (-(-X))$ **by** (*intro* *nat-supermartingale.intro*
submartingale.uminus *nat-submartingale.axioms* *uminus-X.submartingale-nat*) (*auto*
simp add: fun-Compl-def integrable)
thus *?thesis* **unfolding** *fun-Compl-def* **by** *simp*
qed

lemma (*in* *nat-sigma-finite-adapted-process-linorder*) *supermartingale-of-cond-exp-diff-Suc-nonneg*:
assumes *integrable*: $\bigwedge i. \text{integrable } M (X i)$
and $\bigwedge i. AE \xi \text{ in } M. 0 \leq \text{cond-exp } M (F i) (\lambda \xi. X i \xi - X (Suc i) \xi) \xi$
shows *nat-supermartingale* $M F X$
proof (*intro* *supermartingale-nat integrable*)
fix i
show $AE \xi \text{ in } M. X i \xi \geq \text{cond-exp } M (F i) (X (Suc i)) \xi$ **using** *cond-exp-diff[OF*
integrable(1,1), of i i Suc i] *cond-exp-F-meas[OF integrable adapted, of i]* *assms(2)[of*
 $i]$ **by** *fastforce*
qed
end