On the Formalization of Martingales

Ata Keskin

October 31, 2023

Abstract

This thesis presents a formalization of martingales in arbitrary Banach spaces using Isabelle/HOL. We begin by examining formalizations in prominent proof repositories and extend the definition of the conditional expectation operator from real numbers to general Banach spaces, drawing inspiration from prior work. We define filtered measure spaces, adapted, progressively measurable, and predictable processes and rigorously formalize martingales, submartingales, and supermartingales. Additionally, our contributions expand the scope of Bochner integration techniques to general Banach spaces and introduce additional lemmas and induction schemes for integrable functions. Our formalization provides a robust framework for future developments within the theory of stochastic processes. Furthermore, we provide an example demonstrating when a coin toss constitutes a martingale, submartingale, or a supermartingale.

Contents

1	Sup 1.1	Oplementary Lemmas for Measure Spaces Sigma Algebra Generated by a Family of Functions	2	
2	_	pplementary Lemmas for Elementary Metric Spaces Diameter Lemma	3	
3	Sup	oplementary Lemmas for Bochner Integration	5	
	3.1	Integrable Simple Functions	5	
	3.2	Totally Ordered Banach Spaces	14	
	3.3	Integrability and Measurability of the Diameter	16	
	3.4	Auxiliary Lemmas for Set Integrals	18	
	3.5	Averaging Theorem	19	
4	Conditional Expectation in Banach Spaces			
	4.1	Existence	29	
	4.2	Properties	37	
	4.3	Linearly Ordered Banach Spaces	43	
	4.4	Probability Spaces	47	

5	Filte	ered Measure Spaces	53		
	5.1	Filtered Measure	53		
	5.2	Sigma Finite Filtered Measure	54		
	5.3	Finite Filtered Measure	54		
	5.4	Constant Filtration	54		
6	Stochastic Processes 55				
	6.1	Stochastic Process	55		
		6.1.1 Natural Filtration	56		
	6.2	Adapted Process	59		
	6.3	Progressively Measurable Process	62		
	6.4	Predictable Process	65		
7	Martingales 7				
	7.1	Additional Locale Definitions	74		
	7.2	Martingale	76		
	7.3	Submartingale	76		
	7.4	Supermartingale	77		
	7.5	Martingale Lemmas	77		
	7.6	Submartingale Lemmas	80		
	7.7	Supermartingale Lemmas	83		
	7.8	Discrete Time Martingales	87		
	7.9	Discrete Time Submartingales	88		
	7.10	Discrete Time Supermartingales	91		
8	Exa	mple: Coin Toss	92		

```
theory Measure-Space-Supplement
imports HOL-Analysis.Measure-Space
begin
```

1 Supplementary Lemmas for Measure Spaces

1.1 Sigma Algebra Generated by a Family of Functions

```
definition family-vimage-algebra :: 'a set \Rightarrow ('a \Rightarrow 'b) set \Rightarrow 'b measure \Rightarrow 'a
measure where
 family-vimage-algebra \Omega S M \equiv sigma \Omega (\bigcup f \in S. \{f - `A \cap \Omega \mid A. A \in M\})
lemma family-vimage-algebra-singleton: family-vimage-algebra \Omega {f} M = vim-
age-algebra \Omega f M unfolding family-vimage-algebra-def vimage-algebra-def by simp
lemma
 shows sets-family-vimage-algebra: sets (family-vimage-algebra \Omega S M) = sigma-sets
\Omega (\bigcup f \in S. \{f - A \cap \Omega \mid A. A \in M\})
   and space-family-vimage-algebra [simp]: space (family-vimage-algebra \Omega S M) =
\Omega
 by (auto simp add: family-vimage-algebra-def sets-measure-of-conv space-measure-of-conv)
lemma measurable-family-vimage-algebra:
  assumes f \in S f \in \Omega \rightarrow space M
  shows f \in family\text{-}vimage\text{-}algebra\ \Omega\ S\ M \to_M M
  using assms by (intro measurableI, auto simp add: sets-family-vimage-algebra)
{\bf lemma}\ measurable-family-vimage-algebra-singleton:
  assumes f \in \Omega \rightarrow space M
  shows f \in family-vimage-algebra \Omega \{f\} M \to_M M
  using assms measurable-family-vimage-algebra by blast
lemma measurable-family-iff-sets:
  shows (S \subseteq N \to_M M) \longleftrightarrow S \subseteq space N \to space M \land family-vimage-algebra
(space\ N)\ S\ M\subseteq N
proof (standard, goal-cases)
  case 1
  hence subset: S \subseteq space \ N \rightarrow space \ M using measurable-space by fast
  have \{f - A \cap space \mid A \mid A \in M\} \subseteq N \text{ if } f \in S \text{ for } f \text{ using } measur-
able\-iff\-sets[unfolded\ family\-vimage\-algebra\-singleton[symmetric],\ of\ f]\ 1\ subset\ that
by (fastforce simp add: sets-family-vimage-algebra)
 \textbf{then show}~? case~\textbf{unfolding}~sets-family-vimage-algebra~\textbf{using}~sets. sigma-algebra-axioms
by (simp add: subset, intro sigma-algebra.sigma-sets-subset, blast+)
next
  case 2
 hence subset: S \subseteq space N \rightarrow space M by simp
```

show ?case

```
proof (standard, goal-cases)
   case (1 x)
   have family-vimage-algebra (space N) \{x\} M \subseteq N by (metis (no-types, lifting)
1 2 sets-family-vimage-algebra SUP-le-iff sigma-sets-le-sets-iff singletonD)
  thus ?case using measurable-iff-sets[unfolded family-vimage-algebra-singleton[symmetric]]
subset[THEN subsetD, OF 1] by fast
  qed
qed
\mathbf{lemma}\ family\text{-}vimage\text{-}algebra\text{-}diff\text{:}
  shows family-vimage-algebra \Omega S M = sigma \Omega (sets (family-vimage-algebra \Omega
(S-I)\ M) \cup family-vimage-algebra\ \Omega\ (S\cap I)\ M)
 using sets.space-closed space-measure-of-conv
 unfolding family-vimage-algebra-def sets-family-vimage-algebra
 by (intro sigma-eqI, blast, fastforce)
    (intro sigma-sets-eqI, blast, simp add: sets-measure-of-conv split: if-splits,
    meson Diff-subset Sup-subset-mono in-mono inf-sup-ord(1) sigma-sets-subseteq
subset-image-iff, fastforce+)
theory Elementary-Metric-Spaces-Supplement
 imports\ HOL-Analysis. Elementary-Metric-Spaces
begin
```

2 Supplementary Lemmas for Elementary Metric Spaces

2.1 Diameter Lemma

```
lemma diameter-comp-strict-mono:
  fixes s :: nat \Rightarrow 'a :: metric-space
  assumes strict-mono r bounded \{s \mid i \mid i. r \mid n \leq i\}
 shows diameter \{s \ (r \ i) \mid i. \ n \leq i\} \leq diameter \{s \ i \mid i. \ r \ n \leq i\}
proof (rule diameter-subset)
  show \{s\ (r\ i)\ |\ i.\ n\leq i\}\subseteq \{s\ i\ |\ i.\ r\ n\leq i\} using assms(1) monotoneD
strict-mono-mono by fastforce
qed (intro assms(2))
lemma diameter-bounded-bound':
  fixes S :: 'a :: metric\text{-}space set
 assumes S: bdd-above (case-prod dist '(S \times S)) x \in S y \in S
  shows dist\ x\ y \leq diameter\ S
  have (x,y) \in S \times S using S by auto
  then have \textit{dist}\ x\ y \leq (\textit{SUP}\ (x,y) \in \textit{S} \times \textit{S}.\ \textit{dist}\ x\ y) by (\textit{rule}\ \textit{cSUP-upper2} \lceil \textit{OF}\ 
assms(1)]) simp
  with \langle x \in S \rangle show ?thesis by (auto simp: diameter-def)
qed
```

```
lemma bounded-imp-dist-bounded:
   assumes bounded (range s)
   shows bounded ((\lambda(i, j). dist (s i) (s j)) '(\{n..\} \times \{n..\}))
  using bounded-dist-comp[OF bounded-fst bounded-snd, OF bounded-Times(1,1)[OF
assms(1,1)] by (rule bounded-subset, force)
lemma cauchy-iff-diameter-tends-to-zero-and-bounded:
   fixes s :: nat \Rightarrow 'a :: metric-space
  shows Cauchy s \longleftrightarrow ((\lambda n. \ diameter \ \{s \ i \mid i. \ i \geq n\}) \longrightarrow 0 \land bounded \ (range
s))
proof -
   have \{s \mid i \mid i \mid N \leq i\} \neq \{\} for N by blast
   hence diameter-SUP: diameter \{s \mid i \mid i. \ N \leq i\} = (SUP(i, j) \in \{N..\} \times \{N..\}).
dist (s i) (s j)) for N unfolding diameter-def by (auto intro!: arg-cong[of - - Sup])
   show ?thesis
   proof (intro iffI)
      assume asm: Cauchy s
      have \exists N. \forall n \geq N. norm (diameter \{s \ i \ | i. n \leq i\}) < e \ \text{if } e\text{-pos: } e > 0 \ \text{for } e
           obtain N where dist-less: dist (s \ n) \ (s \ m) < (1/2) * e \ if \ n \ge N \ m \ge N
for n m using asm e-pos by (metis Cauchy-def mult-pos-pos zero-less-divide-iff
zero-less-numeral zero-less-one)
             fix r assume r \geq N
            hence dist (s \ n) \ (s \ m) < (1/2) * e \ \text{if} \ n \ge r \ m \ge r \ \text{for} \ n \ m \ \text{using} \ dist-less
that by simp
              hence (SUP\ (i, j) \in \{r...\} \times \{r...\}.\ dist\ (s\ i)\ (s\ j)) \le (1/2) * e by (intro\ intro\ int
cSup-least) fastforce+
             also have \dots < e using e-pos by simp
           finally have diameter \{s \mid i \mid i. \ r \leq i\} < e \text{ using } diameter\text{-}SUP \text{ by } presburger
         moreover have diameter \{s \mid i \mid i. r \leq i\} \geq 0 for r unfolding diameter-SUP
using bounded-imp-dist-bounded OF cauchy-imp-bounded, THEN bounded-imp-bdd-above,
OF \ asm] \ \mathbf{by} \ (intro \ cSup-upper2, \ auto)
          ultimately show ?thesis by auto
         thus (\lambda n. \ diameter \ \{s \ i \ | i. \ n \leq i\}) \longrightarrow 0 \land bounded \ (range \ s) using
cauchy-imp-bounded[OF asm] by (simp add: LIMSEQ-iff)
   next
      assume asm: (\lambda n. \ diameter \{s \ i \mid i. \ n \leq i\}) \longrightarrow 0 \land bounded (range s)
      have \exists N. \forall n \geq N. \forall m \geq N. dist(s n)(s m) < e \text{ if } e\text{-pos: } e > 0 \text{ for } e
            obtain N where diam-less: diameter \{s \ i \ | i. \ r \leq i\} < e \ \text{if} \ r \geq N \ \text{for} \ r
using LIMSEQ-D asm(1) e-pos by fastforce
             fix n m assume n \ge N m \ge N
         hence dist(s n)(s m) < e using cSUP-lessD[OF bounded-imp-dist-bounded]THEN
bounded-imp-bdd-above], OF - diam-less[unfolded diameter-SUP]] asm by auto
          }
```

```
thus ?thesis by blast
qed
then show Cauchy s by (simp add: Cauchy-def)
qed
qed
end

theory Bochner-Integration-Supplement
imports HOL—Analysis.Bochner-Integration HOL—Analysis.Set-Integral Elementary-Metric-Spaces-Supplement
begin
```

3 Supplementary Lemmas for Bochner Integration

3.1 Integrable Simple Functions

```
lemma integrable-implies-simple-function-sequence:
  fixes f :: 'a \Rightarrow 'b :: \{banach, second\text{-}countable\text{-}topology\}
  assumes integrable M f
  obtains s where \bigwedge i. simple-function M (s i)
      and \bigwedge i. emeasure M \{ y \in space M. \ s \ i \ y \neq 0 \} \neq \infty
      and \bigwedge x. \ x \in space \ M \Longrightarrow (\lambda i. \ s \ i \ x) \longrightarrow f \ x
      and \bigwedge x \ i. \ x \in space \ M \Longrightarrow norm \ (s \ i \ x) \le 2 * norm \ (f \ x)
proof-
  have f: f \in borel-measurable M (\int x) norm (f(x) \partial M) < \infty using assms
unfolding integrable-iff-bounded by auto
  obtain s where s: \bigwedge i. simple-function M (s i) \bigwedge x. x \in space M \Longrightarrow (\lambda i. s
i \ x) \longrightarrow f \ x \land i \ x. \ x \in space \ M \Longrightarrow norm \ (s \ i \ x) \le 2 * norm \ (f \ x) \ using
borel-measurable-implies-sequence-metric [OF \ f(1)] unfolding norm-conv-dist by
metis
  {
    \mathbf{fix} i
    have (\int x^+ \cdot x \cdot norm \ (s \ i \ x) \ \partial M) \le (\int x^+ \cdot x \cdot norm \ (f \ x) \cdot x) \ \partial M) using
s by (intro nn-integral-mono, auto)
  also have . . . < \infty using f by (simp add: nn-integral-cmult enreal-mult-less-top
ennreal-mult)
    \textbf{finally have} \ sbi: \ Bochner-Integration. simple-bochner-integrable \ M \ (s \ i) \ \textbf{using}
s by (intro\ simple-bochner-integrable I-bounded)\ auto
     hence emeasure M \{y \in space M. s \ i \ y \neq 0\} \neq \infty by (auto intro: inte-
grable I-simple-bochner-integrable\ simple-bochner-integrable. cases)
 thus ?thesis using that s by blast
qed
lemma simple-function-indicator-representation:
  fixes f ::'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
  assumes f: simple-function M f and x: x \in space M
```

```
(is ? l = ? r)
proof -
    have ?r = (\sum y \in f \text{ 'space } M.
       (if y = f x then indicator (f - `\{y\} \cap space M) x *_R y else 0)) by (auto intro!:
   also have ... = indicator (f - `\{fx\} \cap space M) x *_R fx using assms by (auto
dest: simple-functionD)
    also have \dots = f x using x by (auto simp: indicator-def)
    finally show ?thesis by auto
qed
\mathbf{lemma}\ simple-function-indicator-representation-AE:
    fixes f ::'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
    assumes f: simple-function M f
    shows AE x in M. f x = (\sum y \in f \text{ 'space M. indicator } (f - '\{y\} \cap space M) x
   by (metis (mono-tags, lifting) AE-I2 simple-function-indicator-representation f)
lemmas simple-function-scaleR[intro] = simple-function-compose2[\mathbf{where}\ h=(*_R)]
{\bf lemmas}\ integrable-simple-function = simple-bochner-integrable. intros [\it THEN\ has-bochner-integral-simple-bochner-integrable]
THEN\ integrable.intros]
Induction rule for simple integrable functions.
lemma integrable-simple-function-induct[consumes 2, case-names cong indicator
add, induct set: simple-function]:
    fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
    assumes f: simple-function M f emeasure M \{y \in space M. f y \neq 0\} \neq \infty
   assumes cong: \bigwedge f g. simple-function M f \Longrightarrow emeasure M {y \in space\ M.\ f\ y \neq space\ M.\ f\
\theta \} \neq \infty
                                        \implies simple-function M g \implies emeasure M \{y \in space M. g y \neq
\theta \} \neq \infty
                                      \implies (\bigwedge x. \ x \in space \ M \implies f \ x = g \ x) \implies P \ f \implies P \ g
    assumes indicator: \bigwedge A y. A \in sets M \implies emeasure M A < \infty \implies P (\lambda x.
indicator\ A\ x *_R y)
    assumes add: \bigwedge f g. simple-function M f \Longrightarrow emeasure M {y \in space\ M.\ f\ y \neq space\ M.
\theta \} \neq \infty \Longrightarrow
                                        simple-function M g \Longrightarrow emeasure M \{ y \in space M. g y \neq 0 \} \neq
\infty \Longrightarrow
                                           (\bigwedge z. \ z \in space \ M \Longrightarrow norm \ (f \ z + g \ z) = norm \ (f \ z) + norm
(g z)) \Longrightarrow
                                          P f \Longrightarrow P g \Longrightarrow P (\lambda x. f x + g x)
   shows P f
proof-
   let ?f = \lambda x. (\sum y \in f \text{ 'space } M. \text{ indicat-real } (f - `\{y\} \cap \text{ space } M) \ x *_R y)
  have f-ae-eq: f x = ?f x if x \in space M for x using simple-function-indicator-representation <math>OF
   moreover have emeasure M {y \in space\ M. ?f\ y \neq 0} \neq \infty by (metis (no-types,
lifting) Collect-cong calculation f(2))
```

shows $f x = (\sum y \in f \text{ 'space } M. \text{ indicator } (f - \{y\} \cap \text{space } M) \text{ } x *_R y)$

```
simple-function M (\lambda x. \sum y \in S. indicat-real (f - `\{y\} \cap space M) x
*_R y)
                 emeasure M \{y \in space M. (\sum x \in S. indicat-real (f - `\{x\} \cap space A) \} \}
M) \ y *_R x) \neq 0 \} \neq \infty
                 if S \subseteq f 'space M for S using simple-functionD(1)[OF \ assms(1),
THEN rev-finite-subset, OF that that
  proof (induction rule: finite-induct)
   case empty
    {
      case 1
      then show ?case using indicator[of {}] by force
   next
      case 2
      then show ?case by force
   next
      case 3
      then show ?case by force
   }
  next
   case (insert x F)
   have (f - `\{x\} \cap space M) \subseteq \{y \in space M. f y \neq 0\} if x \neq 0 using that by
   moreover have \{y \in space M. f y \neq 0\} = space M - (f - `\{0\} \cap space M)
by blast
     moreover have space M - (f - `\{0\} \cap space M) \in sets M using sim-
ple-functionD(2)[OF f(1)] by blast
     ultimately have fin-0: emeasure M (f - {x} \cap space M) < \infty if x \neq 0
using that by (metis emeasure-mono f(2) infinity-ennreal-def top.not-eq-extremum
top-unique)
   hence fin-1: emeasure M \{ y \in space \ M. \ indicat\ real \ (f - `\{x\} \cap space \ M) \ y *_R \}
x \neq 0 \neq \infty if x \neq 0 by (metis (mono-tags, lifting) emeasure-mono f(1) indica-
tor\text{-}simps(2)\ linorder\text{-}not\text{-}less\ mem\text{-}Collect\text{-}eq\ scaleR\text{-}eq\text{-}0\text{-}iff\ simple\text{-}function}D(2)
subsetI that)
   have *: (\sum y \in insert \ x \ F. \ indicat-real \ (f - `\{y\} \cap space \ M) \ xa *_R y) = (\sum y \in F.
indicat-real\ (f-`\{y\}\cap space\ M)\ xa*_Ry)+indicat-real\ (f-`\{x\}\cap space\ M)
xa *_R x for xa by (metis (no-types, lifting) Diff-empty Diff-insert0 add.commute
insert.hyps(1) insert.hyps(2) sum.insert-remove)
have **: {y \in space\ M. (\sum x \in insert\ x\ F. indicat\text{-}real\ (f\ -\ `\{x\}\ \cap\ space\ M)\ y **_R\ x) \neq \theta} \subseteq {y \in space\ M. (\sum x \in F. indicat\text{-}real\ (f\ -\ `\{x\}\ \cap\ space\ M)\ y **_R\ x)
\neq 0 \} \cup \{ y \in space \ M. \ indicat\ real \ (f - `\{x\} \cap space \ M) \ y *_R x \neq 0 \} \ unfolding \}
* by fastforce
    {
      case 1
      hence x: x \in f 'space M and F: F \subseteq f 'space M by auto
      show ?case
      proof (cases x = \theta)
       \mathbf{case} \ \mathit{True}
```

moreover have $P(\lambda x. \sum y \in S. indicat-real (f - `\{y\} \cap space M) x *_R y)$

```
then show ?thesis unfolding * using insert(3)[OF\ F] by simp
            next
                case False
                have norm-argument: norm ((\sum y \in F. indicat\text{-real } (f - `\{y\} \cap space M) z)
*_R y) + indicat\text{-real } (f - `\{x\} \cap space M) \ z *_R x) = norm \ (\sum y \in F. \ indicat\text{-real } x \in F. \ indicat\text{-real } x \in F.
(f - `\{y\} \cap space\ M)\ z *_R y) + norm\ (indicat-real\ (f - `\{x\} \cap space\ M)\ z *_R x)
if z: z \in space M for z
                proof (cases f z = x)
                    case True
                    have indicat-real (f - (y) \cap space M) z *_R y = 0 if y \in F for y using
True insert(2) z that 1 unfolding indicator-def by force
                  hence (\sum y \in F. indicat\text{-real } (f - `\{y\} \cap space M) z *_R y) = 0 \text{ by } (meson
sum.neutral)
                    then show ?thesis by force
                next
                    case False
                    then show ?thesis by force
                qed
               show ?thesis using False simple-functionD(2)[OF f(1)] insert(3,5)[OF F]
simple-function-scaleR fin-0 fin-1 by (subst *, subst add, subst simple-function-sum)
(blast intro: norm-argument indicator)+
            qed
        \mathbf{next}
            case 2
            hence x: x \in f 'space M and F: F \subseteq f 'space M by auto
           show ?case
            proof (cases x = \theta)
                case True
                then show ?thesis unfolding * using insert(4)[OF\ F] by simp
            next
           then show ?thesis unfolding * using insert(4)[OFF] simple-functionD(2)[OF
f(1)] by fast
            qed
        \mathbf{next}
            case 3
           hence x: x \in f 'space M and F: F \subseteq f 'space M by auto
            show ?case
            proof (cases x = \theta)
                case True
                then show ?thesis unfolding * using insert(5)[OF\ F] by simp
            next
                case False
                 have emeasure M \{y \in space M. (\sum x \in insert \ x \ F. \ indicat\ real \ (f - `\{x\})\}
\cap space M) y *_R x) \neq 0} \leq emeasure M ({y \in \text{space } M. (\sum x \in F. \text{ indicat-real } (f 
-`\{x\} \cap space\ M)\ y *_R x) \neq 0\} \cup \{y \in space\ M.\ indicat\ real\ (f -`\{x\} \cap space\ M)\}
M) y *_{R} x \neq 0
               \mathbf{using} ** simple-functionD(2)[OF\ insert(4)[OF\ F]]\ simple-functionD(2)[OF\ f]
f(1)] by (intro emeasure-mono, force+)
```

```
also have ... \leq emeasure M \{y \in space M. (\sum x \in F. indicat\text{-real } (f - `\{x\})\}
\cap space M) y *_R x) \neq 0} + emeasure M \{y \in space M. indicat-real (f - `\{x\} \cap space M) \}
space M) y *_R x \neq 0}
                  using simple-functionD(2)[OF\ insert(4)[OF\ F]]\ simple-functionD(2)[OF\ f]
f(1)] by (intro emeasure-subadditive, force+)
                also have ... < \infty using insert(5)[OF\ F]\ fin-1[OF\ False] by (simp\ add:
less-top)
              finally show ?thesis by simp
           qed
       }
   qed
   moreover have simple-function M (\lambda x. \sum y \in f 'space M. indicat-real (f - `\{y\})
\cap space M) x *_R y) using calculation by blast
    moreover have P(\lambda x. \sum y \in f \text{ 'space } M. \text{ indicat-real } (f - \{y\} \cap \text{space } M) \text{ } x
*_R y) using calculation by blast
   ultimately show ?thesis by (intro cong[OF - - f(1,2)], blast, presburger+)
Induction rule for non-negative simple integrable functions
lemma integrable-simple-function-induct-nn[consumes 3, case-names cong indica-
tor add, induct set: simple-function]:
    fixes f: 'a \Rightarrow 'b: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
   assumes f: simple-function M f emeasure M \{y \in space M. f y \neq 0\} \neq \infty \land x.
x \in space \ M \longrightarrow f \ x \ge 0
    assumes cong: \bigwedge f g. simple-function M f \Longrightarrow emeasure M \{y \in space M. f y \in space M. f y
\neq 0 \} \neq \infty \Longrightarrow (\bigwedge x. \ x \in space \ M \Longrightarrow f \ x \geq 0) \Longrightarrow simple-function \ M \ g \Longrightarrow
emeasure M \{ y \in space \ M. \ g \ y \neq 0 \} \neq \infty \Longrightarrow (\bigwedge x. \ x \in space \ M \Longrightarrow g \ x \geq 0 )
\implies (\bigwedge x. \ x \in space \ M \implies f \ x = g \ x) \implies P \ f \implies P \ g
   assumes indicator: \bigwedge A y. y \ge 0 \Longrightarrow A \in sets M \Longrightarrow emeasure M A < \infty \Longrightarrow
P(\lambda x. indicator A x *_R y)
    assumes add: \bigwedge f g. (\bigwedge x. \ x \in space \ M \Longrightarrow f \ x \geq 0) \Longrightarrow simple-function \ M \ f
\implies emeasure M \{ y \in space M. f y \neq 0 \} \neq \infty \implies
                                          (\bigwedge x.\ x\in \mathit{space}\ M\Longrightarrow g\ x\geq 0)\Longrightarrow \mathit{simple-function}\ M\ g\Longrightarrow
emeasure M \{y \in space \ M. \ g \ y \neq 0\} \neq \infty \Longrightarrow (\bigwedge z. \ z \in space \ M \Longrightarrow norm \ (f \ z + g \ z) = norm \ (f \ z) + norm
(g z)) \Longrightarrow
                                       P f \Longrightarrow P g \Longrightarrow P (\lambda x. f x + g x)
   shows Pf
proof-
   let ?f = \lambda x. (\sum y \in f \text{ 'space } M. \text{ indicat-real } (f - `\{y\} \cap \text{space } M) \ x *_R y)
  have f-ae-eq: f x = ?f x if x \in space M for x using simple-function-indicator-representation <math>OF
f(1) that .
   moreover have emeasure M {y \in space\ M. ?f\ y \neq 0} \neq \infty by (metis (no-types,
lifting) Collect-cong calculation f(2))
    moreover have P (\lambda x. \sum y \in S. indicat-real (f - `\{y\} \cap space\ M)\ x *_R y)
                              simple-function M (\lambda x. \sum y \in S. indicat-real (f - `\{y\} \cap space M) x
*_R y)
                               emeasure M {y \in space M. (\sum x \in S. indicat-real (f - `\{x\} \cap space
```

```
M) y *_R x) \neq 0 \} \neq \infty
             \bigwedge x. \ x \in space \ M \Longrightarrow 0 \le (\sum y \in S. \ indicat\ real \ (f - `\{y\} \cap space \ M)
x *_R y
               if S \subseteq f 'space M for S using simple-functionD(1)[OF \ assms(1),
THEN rev-finite-subset, OF that that
 proof (induction rule: finite-subset-induct')
   case empty
    {
     case 1
     then show ?case using indicator[of 0 \ \{\}] by force
   next
     then show ?case by force
   next
     case 3
     then show ?case by force
   next
     case 4
     then show ?case by force
   }
  \mathbf{next}
   case (insert x F)
   have (f - `\{x\} \cap space M) \subseteq \{y \in space M. f y \neq 0\}  if x \neq 0 using that by
   moreover have \{y \in space \ M. \ f \ y \neq \emptyset\} = space \ M - (f - `\{\emptyset\} \cap space \ M)
by blast
    moreover have space M – (f – '\{\theta\} \cap space M) \in sets M using sim-
ple-functionD(2)[OF f(1)] by blast
    ultimately have fin-0: emeasure M (f - {x} \cap space M) < \infty if x \neq 0
using that by (metis emeasure-mono f(2) infinity-enrical-def top.not-eq-extremum
top-unique)
  hence fin-1: emeasure M {y \in space\ M. indicat-real (f - `\{x\} \cap space\ M)\ y *_R
x \neq 0 \neq \infty if x \neq 0 by (metis (mono-tags, lifting) emeasure-mono f(1) indica-
tor\text{-}simps(2)\ linorder\text{-}not\text{-}less\ mem\text{-}Collect\text{-}eq\ scaleR\text{-}eq\text{-}0\text{-}iff\ simple\text{-}function}D(2)
subsetI that)
   have nonneg-x: x \ge 0 using insert f by blast
space M) xa *_R x for xa by (metis (no-types, lifting) add.commute insert.hyps(1)
insert.hyps(4) sum.insert)
   have **: \{y \in space \ M. \ (\sum x \in insert \ x \ F. \ indicat\ real \ (f - `\{x\} \cap space \ M) \ y \}
*_R x) \neq 0 \subseteq \{y \in space \ M. \ (\sum x \in F. \ indicat-real \ (f - `\{x\} \cap space \ M) \ y *_R x)\}
\neq 0} \cup \{y \in space M. indicat-real (f - `\{x\} \cap space M) y *_R x \neq 0\} unfolding
* by fastforce
    {
     case 1
     show ?case
     proof (cases x = \theta)
```

```
case True
               then show ?thesis unfolding * using insert by simp
            next
                case False
                have norm-argument: norm ((\sum y \in F. indicat\text{-real } (f - `\{y\} \cap space M) z)
*_R y) + indicat\text{-real } (f - `\{x\} \cap space M) \ z *_R x) = norm \ (\sum y \in F. \ indicat\text{-real})
(f - `\{y\} \cap space M) \ z *_R y) + norm (indicat-real (f - `\{x\} \cap space M) \ z *_R x)
if z: z \in space M for z
               proof (cases f z = x)
                   case True
                   have indicat-real (f - (y) \cap space M) z *_R y = 0 if y \in F for y using
 True insert z that 1 unfolding indicator-def by force
                 hence (\sum y \in F. indicat\text{-real } (f - `\{y\} \cap space M) \ z *_R y) = 0 \text{ by } (meson
sum.neutral)
                   thus ?thesis by force
               qed (force)
             show ?thesis using False fin-0 fin-1 f norm-argument by (subst *, subst add,
presburger add: insert, intro insert, intro insert, fastforce simp add: indicator-def
intro!: insert(2) f(3), auto intro!: indicator insert nonneg-x)
            qed
       next
            case 2
            show ?case
            proof (cases x = \theta)
               {\bf case}\ {\it True}
               then show ?thesis unfolding * using insert by simp
            next
               case False
              then show ?thesis unfolding * using insert simple-functionD(2)[OF f(1)]
by fast
            qed
       next
            case 3
           show ?case
            proof (cases x = \theta)
               case True
               then show ?thesis unfolding * using insert by simp
            next
                case False
                 have emeasure M \{y \in space M. (\sum x \in insert \ x \ F. \ indicat-real \ (f - `\{x\})\}
\cap space M) y *_R x) \neq 0} \leq emeasure M ({y \in \text{space } M. (\sum x \in F. \text{ indicat-real } (f 
-`\{x\} \cap space\ M)\ y *_R x) \neq 0\} \cup \{y \in space\ M.\ indicat\ real\ (f -`\{x\} \cap space\ M)\}
M) \ y *_{R} x \neq 0 \})
                 using ** simple-functionD(2)[OF\ insert(6)]\ simple-functionD(2)[OF\ f(1)]
insert.IH(2) by (intro emeasure-mono, blast, simp)
               also have ... \leq emeasure M \{y \in space M. (\sum x \in F. indicat\text{-real } (f - `\{x\})\}
\cap space M) y *_R x) \neq 0} + emeasure M \{y \in \text{space } M \text{. indicat-real } (f - `\{x\} \cap Y) \}
space M) y *_R x \neq 0}
                      using simple-functionD(2)[OF\ insert(6)]\ simple-functionD(2)[OF\ f(1)]
```

```
by (intro emeasure-subadditive, force+)
       also have ... < \infty using insert(7) fin-1[OF False] by (simp add: less-top)
       finally show ?thesis by simp
     qed
   next
     case (4 \xi)
    thus ? case using insert nonneg-x f(3) by (auto simp add: scaleR-nonneg-nonneg
intro: sum-nonneg)
   }
 \mathbf{qed}
 moreover have simple-function M (\lambda x. \sum y \in f 'space M. indicat-real (f - `\{y\})
\cap space M) x *_R y) using calculation by blast
  moreover have P(\lambda x. \sum y \in f \text{ 'space } M. \text{ indicat-real } (f - \{y\} \cap \text{space } M) \text{ } x
*_R y) using calculation by blast
 moreover have \bigwedge x. x \in space M \Longrightarrow 0 \le f x using f(3) by simp
  ultimately show ?thesis by (intro cong[OF - - - f(1,2)], blast, blast, fast)
presburger +
qed
lemma finite-nn-integral-imp-ae-finite:
  fixes f :: 'a \Rightarrow ennreal
  assumes f \in borel-measurable M (\int x. f x \partial M) < \infty
  shows AE \ x \ in \ M. \ f \ x < \infty
proof (rule ccontr, goal-cases)
  case 1
  let ?A = space M \cap \{x. f x = \infty\}
  have *: emeasure M?A > 0 using 1 assms(1) by (metis (mono-tags, lifting)
assms(2) eventually-mono infinity-enreal-def nn-integral-noteq-infinite top.not-eq-extremum)
  have (\int_{-\infty}^{+\infty} x \cdot f \cdot x * indicator ?A \times \partial M) = (\int_{-\infty}^{+\infty} x \cdot \infty * indicator ?A \times \partial M) by
(metis (mono-tags, lifting) indicator-inter-arith indicator-simps(2) mem-Collect-eq
mult-eq-\theta-iff)
 also have \dots = \infty * emeasure M ?A using assms(1) by (intro nn-integral-cmult-indicator,
simp)
  also have \dots = \infty using * by fastforce
  finally have (\int x \cdot f \cdot x \cdot indicator ?A \cdot x \cdot \partial M) = \infty.
  moreover have (\int x \cdot f \cdot x \cdot indicator ?A \cdot x \cdot \partial M) \leq (\int x \cdot f \cdot x \cdot \partial M) by (intro
nn	ext{-}integral	ext{-}mono, \ simp \ add: \ indicator	ext{-}def)
  ultimately show ?case using assms(2) by simp
qed
Convergence in L1-Norm implies existence of a subsequence which conver-
gences almost everywhere. This lemma is easier to use than the existing one
in HOL-Analysis. Bochner-Integration
lemma cauchy-L1-AE-cauchy-subseq:
  fixes s:: nat \Rightarrow 'a \Rightarrow 'b::\{banach, second-countable-topology\}
  assumes [measurable]: \bigwedge n. integrable \ M \ (s \ n)
     and \bigwedge e. \ e > 0 \Longrightarrow \exists N. \ \forall i \geq N. \ \forall j \geq N. \ LINT \ x | M. \ norm \ (s \ i \ x - s \ j \ x) < e
  obtains r where strict-mono r AE x in M. Cauchy (\lambda i.\ s\ (r\ i)\ x)
proof-
```

```
have \exists r. \forall n. (\forall i \geq r \ n. \ \forall j \geq r \ n. \ LINT \ x | M. \ norm \ (s \ i \ x - s \ j \ x) < (1 \ / \ 2) 
n) \wedge (r (Suc \ n) > r \ n)
 proof (intro dependent-nat-choice, goal-cases)
   case 1
   then show ?case using assms(2) by presburger
 next
   case (2 x n)
   obtain N where *: LINT x|M. norm (s i x - s j x) < (1 / 2) \cap Suc n if i \ge
N j \geq N for i j using assms(2)[of (1 / 2) \cap Suc n] by auto
     fix i j assume i \geq max \ N \ (Suc \ x) \ j \geq max \ N \ (Suc \ x)
     hence integral<sup>L</sup> M (\lambda x. norm (s i x - s j x)) < (1 / 2) \hat{S}uc n using * by
fast force
   then show ?case by fastforce
 qed
 then obtain r where strict-mono: strict-mono r and \forall i \geq r \ n. \ \forall j \geq r \ n. \ LINT
x|M. norm (s \ i \ x - s \ j \ x) < (1 \ / \ 2) \ \hat{} \ n for n using strict-mono-Suc-iff by blast
 hence r-is: LINT x|M. norm (s(r(Suc n)) x - s(r n) x) < (1/2) \cap n for n
by (simp add: strict-mono-leD)
  define g where g = (\lambda n \ x. \ (\sum i \le n. \ ennreal \ (norm \ (s \ (r \ (Suc \ i)) \ x - s \ (r \ i))))
  define g' where g' = (\lambda x. \sum i. ennreal (norm (s (r (Suc i)) x - s (r i) x)))
 have integrable-g: (\int + x. g \ n \ x \ \partial M) < 2 \ \text{for} \ n
   have (\int x. g \, n \, x \, \partial M) = (\int x. (\sum i \leq n. ennreal (norm (s (r (Suc i)) x - i)))
s\ (r\ i)\ x)))\ \partial M) using g\text{-}def by simp
   also have ... = (\sum i \le n. (\int + x. ennreal (norm (s (r (Suc i)) x - s (r i) x)))
\partial M)) by (intro\ nn\mathchar`-integral\mathchar`-sum,\ simp)
     also have ... = (\sum i \le n. LINT x|M. norm (s (r (Suc i)) x - s (r i) x))
unfolding dist-norm using assms(1) by (subst nn-integral-eq-integral[OF inte-
grable-norm], auto)
  also have ... < ennreal (\sum i \le n. (1/2)^{i}) by (intro ennreal-lessI[OF sum-pos
sum-strict-mono[OF finite-atMost - r-is]], auto)
   also have ... \leq ennreal\ 2 unfolding sum-gp0[of\ 1\ /\ 2\ n] by (intro ennreal-leI,
auto)
   finally show (\int + x. g n x \partial M) < 2 by simp
 have integrable-g': (\int + x. g' x \partial M) \leq 2
    have incseq (\lambda n. \ g \ n \ x) for x by (intro incseq-SucI, auto simp add: g-def
ennreal-leI)
    hence convergent (\lambda n. \ g \ n \ x) for x unfolding convergent-def using LIM-
SEQ-incseq-SUP by fast
     hence (\lambda n. \ g \ n \ x) \longrightarrow g' \ x for x unfolding g-def g'-def by (intro
summable-iff-convergent' [THEN iffD2, THEN summable-LIMSEQ'], blast)
```

```
hence (\int x \cdot g' x \, \partial M) = (\int x \cdot liminf(\lambda n \cdot g \cdot n \cdot x) \, \partial M) by (metis lim-imp-Liminf
trivial-limit-sequentially)
   also have ... \leq liminf(\lambda n. \int + x. g n x \partial M) by (intro nn-integral-liminf, simp
add: g\text{-}def)
   also have ... \leq liminf(\lambda n. 2) using integrable-g by (intro Liminf-mono) (simp
add: order-le-less)
   also have ... = 2 using sequentially-bot tendsto-iff-Liminf-eq-Limsup by blast
   finally show ?thesis.
  qed
 hence AE x in M. g' x < \infty by (intro finite-nn-integral-imp-ae-finite) (auto simp
add: order-le-less-trans g'-def)
  moreover have summable (\lambda i. norm (s (r (Suc i)) x - s (r i) x)) if g' x \neq \infty
for x using that unfolding g'-def by (intro summable-suminf-not-top) fastforce+
  ultimately have ae-summable: AE x in M. summable (\lambda i.\ s\ (r\ (Suc\ i))\ x-s
(r i) x) using summable-norm-cancel unfolding dist-norm by force
   fix x assume summable (\lambda i. s (r (Suc i)) x - s (r i) x)
   hence (\lambda n. \sum i < n. \ s \ (r \ (Suc \ i)) \ x - s \ (r \ i) \ x) \longrightarrow (\sum i. \ s \ (r \ (Suc \ i)) \ x - s \ (r \ i) \ x)
s\ (r\ i)\ x)\ \mathbf{using}\ summable\text{-}LIMSEQ\ \mathbf{by}\ blast
   moreover have (\lambda n. (\sum i < n. s (r (Suc i)) x - s (r i) x)) = (\lambda n. s (r n) x - s (r i) x)
s\ (r\ \theta)\ x) using sum-less Than-telescope by fast force
   ultimately have (\lambda n. \ s \ (r \ n) \ x - s \ (r \ 0) \ x) \longrightarrow (\sum i. \ s \ (r \ (Suc \ i)) \ x - s
(r \ i) \ x) by argo
   hence (\lambda n. \ s \ (r \ n) \ x - s \ (r \ 0) \ x + s \ (r \ 0) \ x) \longrightarrow (\sum i. \ s \ (r \ (Suc \ i)) \ x - s
(r \ i) \ x) + s \ (r \ \theta) \ x \ by \ (intro \ isCont-tendsto-compose[of - \lambda z. \ z + s \ (r \ \theta) \ x], \ auto)
   hence Cauchy (\lambda n. \ s \ (r \ n) \ x) by (simp \ add: LIMSEQ-imp-Cauchy)
 hence AE x in M. Cauchy (\lambda i.\ s\ (r\ i)\ x) using ae-summable by fast
  thus ?thesis by (rule\ that[OF\ strict-mono(1)])
qed
3.2
        Totally Ordered Banach Spaces
lemma integrable-max[simp, intro]:
 fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology}\}
  assumes fg[measurable]: integrable M f integrable M g
  shows integrable M (\lambda x. max (f x) (g x))
proof (rule Bochner-Integration.integrable-bound)
  {
   fix x y :: 'b
   have norm (max \ x \ y) \le max \ (norm \ x) \ (norm \ y) by linarith
   also have ... \leq norm \ x + norm \ y \ by \ simp
   finally have norm (max \ x \ y) \leq norm \ (norm \ x + norm \ y) by simp
```

thus $AE \ x \ in \ M. \ norm \ (max \ (f \ x) \ (g \ x)) \leq norm \ (norm \ (f \ x) + norm \ (g \ x))$ by

qed (auto intro: Bochner-Integration.integrable-add[OF integrable-norm[OF fq(1)]

```
lemma integrable-min[simp, intro]:
  fixes f: 'a \Rightarrow 'b: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology}\}
  assumes [measurable]: integrable M f integrable M q
  shows integrable M (\lambda x. min (f x) (g x))
proof -
  have norm (min (f x) (q x)) \le norm (f x) \lor norm (min (f x) (q x)) \le norm (q x)
x) for x by linarith
 thus ?thesis by (intro integrable-bound[OF integrable-max[OF integrable-norm(1,1),
OF \ assms]], \ auto)
qed
lemma integral-nonneg-AE-banach:
  fixes f::'a \Rightarrow 'b:: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or\text{-}
dered-real-vector}
 assumes [measurable]: f \in borel-measurable M and nonneg: AE x in M. 0 \le f x
  shows 0 \leq integral^L M f
proof cases
  assume integrable: integrable M f
  hence max: (\lambda x. \ max \ \theta \ (f \ x)) \in borel-measurable M \ \land x. \ \theta \leq max \ \theta \ (f \ x)
integrable M (\lambda x. max \theta (f x)) by auto
  hence 0 \leq integral^L M(\lambda x. max \theta(f x))
  proof -
  obtain s where *: \bigwedge i. simple-function M (s i)
                    \bigwedge i. emeasure M \{ y \in space M. \ s \ i \ y \neq 0 \} \neq \infty
                    \bigwedge x. \ x \in space \ M \Longrightarrow (\lambda i. \ s \ i \ x) \longrightarrow f \ x
                      \bigwedge x \ i. \ x \in space \ M \Longrightarrow norm \ (s \ i \ x) \le 2 * norm \ (f \ x)  using
integrable-implies-simple-function-sequence[OF integrable] by blast
    have simple: \bigwedge i. simple-function M (\lambda x. max \theta (s i x)) using * by fast
    have \bigwedge i. \{y \in space \ M. \ max \ 0 \ (s \ i \ y) \neq 0\} \subseteq \{y \in space \ M. \ s \ i \ y \neq 0\}
unfolding max-def by force
   moreover have \bigwedge i. \{y \in space \ M. \ s \ i \ y \neq 0\} \in sets \ M \ using * by \ measurable
     ultimately have \bigwedge i. emeasure M \{y \in space M. max 0 (s i y) \neq 0\} \leq
emeasure M \{ y \in space \ M. \ s \ i \ y \neq 0 \} by (rule emeasure-mono)
    hence **:\land i. emeasure M \{ y \in space M. max \theta (s i y) \neq \theta \} \neq \infty using *(2)
by (auto intro: order.strict-trans1 simp add: top.not-eq-extremum)
    have \bigwedge x. \ x \in space \ M \Longrightarrow (\lambda i. \ max \ \theta \ (s \ i \ x)) \longrightarrow max \ \theta \ (f \ x) \ using *(3)
tendsto-max by blast
    moreover have \bigwedge x \ i. \ x \in space \ M \Longrightarrow norm \ (max \ 0 \ (s \ i \ x)) \leq norm \ (2 *_R)
f(x) using *(4) unfolding max-def by auto
   ultimately have tendsto: (\lambda i. integral^L M (\lambda x. max \theta (s i x))) \longrightarrow integral^L
M (\lambda x. max \theta (f x))
         using borel-measurable-simple-function simple integrable by (intro inte-
gral-dominated-convergence[OF\ max(1)\ -\ integrable-norm[OF\ integrable-scaleR-right],
of - 2f], blast+)
      fix h: 'a \Rightarrow 'b: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
```

 $integrable-norm[OF\ fg(2)]])$

dered-real-vector}

```
assume simple-function M h emeasure M \{y \in space M. h y \neq 0\} \neq \infty \land x.
x \in space \ M \longrightarrow h \ x \geq \, \theta
     hence *: integral^L M h \ge 0
     proof (induct rule: integrable-simple-function-induct-nn)
       case (conq f q)
       then show ?case using Bochner-Integration.integral-cong by force
     next
       case (indicator A y)
       hence A \neq \{\} \Longrightarrow y \geq \theta using sets.sets-into-space by fastforce
          then show ?case using indicator by (cases A = \{\}), auto simp add:
scaleR-nonneg-nonneg)
     next
       case (add f g)
       then show ?case by (simp add: integrable-simple-function)
   thus ?thesis using LIMSEQ-le-const[OF tendsto, of 0] ** simple by fastforce
  qed
 also have ... = integral^L M f using nonneg by (auto intro: integral-cong-AE)
  finally show ?thesis.
qed (simp add: not-integrable-integral-eq)
lemma integral-mono-AE-banach:
  fixes f g :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes integrable M f integrable M g AE x in M. f x \leq g x
 shows integral^L M f \leq integral^L M g
 using integral-nonneg-AE-banach of \lambda x. gx - fx assms Bochner-Integration.integral-diff OF
assms(1,2)] by force
lemma integral-mono-banach:
  fixes fg: 'a \Rightarrow 'b: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes integrable M f integrable M g \land x. x \in space M \Longrightarrow f x \leq g x
 shows integral^L M f \leq integral^L M g
 using integral-mono-AE-banach assms by blast
3.3
       Integrability and Measurability of the Diameter
context
 fixes s:: nat \Rightarrow 'a \Rightarrow 'b:: \{second\text{-}countable\text{-}topology, banach} \} and M
 assumes bounded: \bigwedge x. \ x \in space \ M \Longrightarrow bounded \ (range \ (\lambda i. \ s \ i \ x))
begin
lemma borel-measurable-diameter:
 assumes [measurable]: \bigwedge i. (s i) \in borel-measurable M
 shows (\lambda x. \ diameter \{s \ i \ x \ | i. \ n \leq i\}) \in borel-measurable M
proof -
 have \{s \ i \ x \mid i.\ N \leq i\} \neq \{\} for x \ N by blast
```

```
hence diameter-SUP: diameter \{s \ i \ x \ | i.\ N \leq i\} = (SUP\ (i,j) \in \{N..\} \times \{N..\}.
dist\ (s\ i\ x)\ (s\ j\ x)) for x\ N unfolding diameter-def by (auto\ intro!:\ arg-cong[of\ -
- Sup
   have case-prod dist '(\{s \ i \ x \ | i. \ n \leq i\} \times \{s \ i \ x \ | i. \ n \leq i\}) = ((\lambda(i, j). \ dist \ (s \ i \ x \ | i. \ n \leq i))
(s \ j \ x) \ (\{n..\} \times \{n..\})) for x \ by \ fast
    hence *: (\lambda x. \ diameter \{s \ i \ x \ | i. \ n \le i\}) = (\lambda x. \ Sup \ ((\lambda(i, j). \ dist \ (s \ i \ x) \ (s \ j)) \}
x)) '(\{n..\} \times \{n..\}))) using diameter-SUP by (simp add: case-prod-beta')
   have bounded ((\lambda(i, j). \ dist \ (s \ i \ x) \ (s \ j \ x)) \ `(\{n..\} \times \{n..\})) \ \text{if} \ x \in space \ M \ \text{for}
x by (rule bounded-imp-dist-bounded[OF bounded, OF that])
   hence bdd: bdd-above ((\lambda(i,j), dist (s i x) (s j x)) `(\{n..\} \times \{n..\})) if x \in space
M for x using that bounded-imp-bdd-above by presburger
   have fst p \in borel-measurable M snd p \in borel-measurable M if p \in s '\{n..\} \times
s ` \{n..\}  for p using that by fastforce+
   hence (\lambda x. \ fst \ p \ x - snd \ p \ x) \in borel-measurable M \ \textbf{if} \ p \in s \ `\{n..\} \times s \ `\{n..\}
for p using that borel-measurable-diff by simp
    hence (\lambda x. \ case \ p \ of \ (f, \ g) \Rightarrow \ dist \ (f \ x) \ (g \ x)) \in borel-measurable \ M \ \textbf{if} \ p \in s
\{n..\} \times s \ (n..\}  for p unfolding dist-norm using that by measurable
     moreover have countable (s '\{n..\} × s '\{n..\}) by (intro countable-SIGMA
countable-image, auto)
     ultimately show ?thesis unfolding * by (auto intro!: borel-measurable-cSUP
bdd)
qed
{\bf lemma}\ integrable	ext{-}bound	ext{-}diameter:
    fixes f :: 'a \Rightarrow real
    assumes integrable M f
            and [measurable]: \land i. (s i) \in borel-measurable M
            and \bigwedge x \ i. \ x \in space \ M \Longrightarrow norm \ (s \ i \ x) \le f \ x
        shows integrable M (\lambda x. diameter \{s \ i \ x \ | i. \ n \leq i\})
proof -
    have \{s \ i \ x \mid i.\ N \leq i\} \neq \{\} for x \ N by blast
   hence diameter-SUP: diameter \{s \ i \ x \ | i.\ N \leq i\} = (SUP\ (i,j) \in \{N..\} \times \{N..\}.
dist (s i x) (s j x)) for x N unfolding diameter-def by (auto intro!: arg-cong[of -
- Sup
    {
        fix x assume x: x \in space M
        let S = (\lambda(i, j)). dist (s \ i \ x) \ (s \ j \ x)) '(\{n..\} \times \{n..\})
        have case-prod dist '(\{s \ i \ x \ | i. \ n \leq i\}) \times \{s \ i \ x \ | i. \ n \leq i\}) = (\lambda(i, j). \ dist \ (s \ i \ j) = (\lambda(i, j)) + (\lambda(i, j))
(s \ j \ x)) '(\{n..\} \times \{n..\}) by fast
        hence *: diameter \{s \ i \ x \ | i. \ n \leq i\} = Sup \ ?S \ using \ diameter-SUP \ by \ (simp)
add: case-prod-beta')
        have bounded ?S by (rule bounded-imp-dist-bounded[OF bounded[OF x]])
      hence Sup-S-nonneg: 0 \le Sup ?S by (auto intro!: cSup-upper2 x bounded-upp-bdd-above)
          have dist (s \ i \ x) \ (s \ j \ x) \le 2 * f \ x for i \ j \ by \ (intro \ dist-triangle 2 \ | THEN
```

order-trans, of - 0]) (metis norm-conv-dist assms(3) x add-mono mult-2)

```
hence \forall c \in ?S. \ c \leq 2 * f x \text{ by } force
      hence Sup ?S \le 2 * fx by (intro cSup-least, auto)
       hence norm\ (Sup\ ?S) \le 2*norm\ (f\ x) using Sup\text{-}S\text{-}nonneg by auto
       also have ... = norm (2 *_R f x) by simp
       finally have norm (diameter \{s \ i \ x \ | i. \ n \leq i\}) \leq norm \ (2 *_R f x) unfolding
  hence AE \ x \ in \ M. \ norm \ (diameter \ \{s \ i \ x \ | i. \ n \leq i\}) \leq norm \ (2 *_R f \ x) \ \mathbf{by} \ blast
  thus integrable M (\lambda x. diameter \{s \ i \ x \ | i. \ n \leq i\}) using borel-measurable-diameter
\textbf{by } (intro\ Bochner-Integration.integrable-bound [OF\ assms(1)] THEN\ integrable-scaleR-right [of\ Assms(1
2]]], measurable)
qed
end
               Auxiliary Lemmas for Set Integrals
3.4
lemma set-integral-scaleR-left:
   assumes A \in sets \ M \ c \neq 0 \Longrightarrow integrable \ M \ f
   shows LINT t:A|M. f t *_R c = (LINT t:A|M. f t) *_R c
    unfolding set-lebesgue-integral-def
    using integrable-mult-indicator[OF assms]
   by (subst integral-scaleR-left[symmetric], auto)
lemma nn-set-integral-eq-set-integral:
    assumes [measurable]: integrable M f
          and AE x \in A in M. 0 \le f x A \in sets M
       shows (\int x \in A. f x \partial M) = (\int x \in A. f x \partial M)
proof-
   have (\int x \cdot indicator A \cdot x *_R f \cdot x \cdot \partial M) = (\int x \in A \cdot f \cdot x \cdot \partial M)
  unfolding set-lebesque-integral-def using assms(2) by (intro nn-integral-eq-integral of
- \lambda x. indicat-real A \times_R f x, blast intro: assms integrable-mult-indicator, fastforce)
  moreover have (\int_{-\infty}^{+\infty} x \cdot indicator A \times_R f \times \partial M) = (\int_{-\infty}^{+\infty} x \in A \cdot f \times \partial M) by (metis
ennreal-0 indicator-simps(1) indicator-simps(2) mult.commute mult-1 mult-zero-left
real-scaleR-def)
   ultimately show ?thesis by argo
qed
lemma set-integral-restrict-space:
   fixes f :: 'a \Rightarrow 'b :: \{banach, second-countable-topology\}
   assumes \Omega \cap space M \in sets M
   shows set-lebesque-integral (restrict-space M \Omega) A f = set-lebesque-integral M A
(\lambda x. indicator \Omega x *_R f x)
    unfolding set-lebesgue-integral-def
  by (subst integral-restrict-space, auto intro!: integrable-mult-indicator assms simp:
mult.commute)
\mathbf{lemma}\ \mathit{set-integral-const}:
   fixes c :: 'b::\{banach, second-countable-topology\}
   assumes A \in sets \ M \ emeasure \ M \ A \neq \infty
```

```
shows set-lebesgue-integral M A (\lambda-. c) = measure M A *_R c
 unfolding set-lebesgue-integral-def
 using assms by (metis has-bochner-integral-indicator has-bochner-integral-integral-eq
infinity-ennreal-def less-top)
lemma set-integral-mono-banach:
  fixes f g :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
  assumes set-integrable M A f set-integrable M A g
   \bigwedge x. \ x \in A \Longrightarrow f \ x \leq g \ x
 shows (LINT x:A|M. fx) \le (LINT x:A|M. gx)
 using assms unfolding set-integrable-def set-lebesgue-integral-def
 by (auto intro: integral-mono-banach split: split-indicator)
lemma set-integral-mono-AE-banach:
  fixes f g :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes set-integrable M A f set-integrable M A g AE x \in A in M. f x \leq g x
  shows set-lebesgue-integral M A f \leq set-lebesgue-integral M A g using assms
unfolding set-lebesque-integral-def by (auto simp add: set-integrable-def intro!:
integral-mono-AE-banach[of\ M\ \lambda x.\ indicator\ A\ x*_R fx\ \lambda x.\ indicator\ A\ x*_R g\ x],
simp add: indicator-def)
        Averaging Theorem
3.5
{f lemma}\ balls	ext{-}countable	ext{-}basis:
  obtains D :: 'a :: \{metric\text{-}space, second\text{-}countable\text{-}topology}\} set
  where topological-basis (case-prod ball '(D \times (\mathbb{Q} \cap \{0 < ... \})))
   and countable D
   and D \neq \{\}
proof -
 obtain D: 'a set where dense-subset: countable D D \neq \{\} [open U; U \neq \{\}]
\implies \exists y \in D. \ y \in U \ \text{for} \ U \ \text{using } countable\text{-}dense\text{-}exists \ \text{by } blast
 have topological-basis (case-prod ball '(D \times (\mathbb{Q} \cap \{0 < ... \})))
 proof (intro topological-basis-iff[THEN iffD2], fast, clarify)
   fix U and x :: 'a assume asm: open U x \in U
   obtain e where e: e > 0 ball x \in U using asm openE by blast
  obtain y where y: y \in D y \in ball x (e / 3) using dense-subset(3)[OF open-ball,
of x \in /3 centre-in-ball [THEN iffD2, OF divide-pos-pos[OF e(1), of 3]] by force
  obtain r where r: r \in \mathbb{Q} \cap \{e/3 < ... < e/2\} unfolding Rats-def using of-rat-dense[OF]
divide-strict-left-mono[OF - e(1)], of 2 3 by auto
   have *: x \in ball\ y\ r\ using\ r\ y\ by\ (simp\ add:\ dist-commute)
   hence ball y r \subseteq U using r by (intro order-trans[OF - e(2)], simp, metric)
    moreover have ball y \in (case-prod\ ball\ `(D \times (\mathbb{Q} \cap \{0 < ..\}))) using y(1)
r by force
   ultimately show \exists B' \in (case\text{-prod ball } (D \times (\mathbb{Q} \cap \{0 < ..\}))). \ x \in B' \wedge B' \subseteq
U using * by meson
```

qed

```
thus ?thesis using that dense-subset by blast
qed
context sigma-finite-measure
begin
lemma sigma-finite-measure-induct[case-names finite-measure, consumes 0]:
 assumes \bigwedge(N :: 'a \ measure) \ \Omega. finite-measure N
                            \implies N = restrict\text{-}space \ M \ \Omega
                            \Longrightarrow \Omega \in \mathit{sets}\ M
                            \implies emeasure \ N \ \Omega \neq \infty
                            \implies emeasure \ N \ \Omega \neq 0
                            \implies almost\text{-}everywhere \ N \ Q
     and [measurable]: Measurable.pred\ M\ Q
 shows almost-everywhere M Q
proof -
  have *: almost-everywhere N Q if finite-measure N N = restrict-space M \Omega \Omega
\in sets M emeasure N \Omega \neq \infty for N \Omega using that by (cases emeasure N \Omega = 0,
auto intro: emeasure-0-AE \ assms(1))
 obtain A :: nat \Rightarrow 'a \text{ set where } A : range A \subseteq sets M (\bigcup i. A i) = space M \text{ and}
emeasure-finite: emeasure M (A i) \neq \infty for i using sigma-finite by metis
 note A(1)[measurable]
 have space-restr: space (restrict-space M(A i)) = A i for i unfolding space-restrict-space
\mathbf{by} \ simp
  {
   \mathbf{fix} i
   have *: \{x \in A \ i \cap space \ M. \ Q \ x\} = \{x \in space \ M. \ Q \ x\} \cap (A \ i) by fast
  have Measurable.pred (restrict-space M(A i)) Q using A by (intro measurable I,
auto simp add: space-restr intro!: sets-restrict-space-iff[THEN iffD2], measurable,
auto)
 }
 note this[measurable]
   \mathbf{fix} i
    have finite-measure (restrict-space M (A i)) using emeasure-finite by (intro
finite-measureI, subst space-restr, subst emeasure-restrict-space, auto)
    hence emeasure (restrict-space M (A i)) \{x \in A \ i. \ \neg Q \ x\} = 0 using emea-
sure-finite by (intro AE-iff-measurable THEN iffD1, OF - - *], measurable, subst
space-restr[symmetric], intro sets.top, auto simp add: emeasure-restrict-space)
  hence emeasure M \{x \in A \ i. \ \neg Q \ x\} = 0 by (subst emeasure-restrict-space[symmetric],
auto)
  }
 hence emeasure M (\bigcup i. \{x \in A \ i : \neg Q \ x\}) = 0 by (intro emeasure-UN-eq-0,
 moreover have (\bigcup i. \{x \in A \ i. \neg Q \ x\}) = \{x \in space \ M. \neg Q \ x\} \text{ using } A \text{ by }
  ultimately show ?thesis by (intro AE-iff-measurable[THEN iffD2], auto)
qed
```

The following lemma allows us to make statements about the behavious of a function almost everywhere, depending on the value it takes on average.

lemma averaging-theorem:

and closed: closed S

assumes [measurable]: integrable M f

fixes $f::- \Rightarrow 'b::\{second\text{-}countable\text{-}topology, banach}\}$

```
and \bigwedge A. A \in sets \ M \Longrightarrow measure \ M \ A > 0 \Longrightarrow (1 \ / measure \ M \ A) *_R
set-lebesgue-integral M A f \in S
   shows AE x in M. f x \in S
proof (induct rule: sigma-finite-measure-induct)
  case (finite-measure N \Omega)
 interpret finite-measure N by (rule finite-measure)
 have integrable[measurable]: integrable N f using assms finite-measure by (auto
simp: integrable-restrict-space integrable-mult-indicator)
  have average: (1 / Sigma-Algebra.measure N A) *<sub>R</sub> set-lebesgue-integral N A f
\in S \text{ if } A \in sets \ N \ measure \ N \ A > 0 \ \text{for } A
 proof -
  have *: A \in sets M using that by (simp add: sets-restrict-space-iff finite-measure)
   have A = A \cap \Omega by (metis finite-measure(2,3) inf.orderE sets.sets-into-space
space-restrict-space\ that(1))
    hence set-lebesgue-integral N A f = set-lebesgue-integral M A f unfolding
finite-measure by (subst set-integral-restrict-space, auto simp add: finite-measure
set-lebesgue-integral-def indicator-inter-arith[symmetric])
    moreover have measure N A = measure M A using that by (auto intro!:
measure-restrict-space simp add: finite-measure sets-restrict-space-iff)
   ultimately show ?thesis using that * assms(3) by presburger
 qed
 obtain D: 'b set where balls-basis: topological-basis (case-prod ball '(D \times (\mathbb{Q}
\cap \{0 < ... \})) and countable-D: countable D using balls-countable-basis by blast
  have countable-balls: countable (case-prod ball ' (D \times (\mathbb{Q} \cap \{\theta < ...\}))) using
countable-rat countable-D by blast
  obtain B where B-balls: B \subseteq case\text{-prod ball} \ (D \times (\mathbb{Q} \cap \{0 < ..\})) \cup B = -S
using topological-basis[THEN iffD1, OF balls-basis] open-Compl[OF assms(2)] by
meson
 hence countable-B: countable B using countable-balls countable-subset by fast
 define b where b = from\text{-}nat\text{-}into\ (B \cup \{\{\}\}\})
 have B \cup \{\{\}\} \neq \{\} by simp
 have range-b: range b = B \cup \{\{\}\} using countable-B by (auto simp add: b-def
intro!: range-from-nat-into)
 have open-b: open (b i) for i unfolding b-def using B-balls open-ball from-nat-into[of
B \cup \{\{\}\}\ i by force
 have Union-range-b: \bigcup (range\ b) = -S using B-balls range-b by simp
  {
```

```
fix v r assume ball-in-Compl: ball v r \subseteq -S
   define A where A = f - `ball v r \cap space N
   have dist-less: dist (f x) v < r if x \in A for x using that unfolding A-def
vimage-def by (simp add: dist-commute)
    hence AE-less: AE x \in A in N. norm (f x - v) < r by (auto simp add:
dist-norm)
   have *: A \in sets \ N  unfolding A-def by simp
   have emeasure NA = 0
   proof -
    {
      assume asm: emeasure NA > 0
      hence measure-pos: measure NA > 0 unfolding emeasure-eq-measure by
simp
    A) *_R set-lebesque-integral N A (\lambda x. f x - v) using integrable integrable-const * by
(subst\ set\ -integral\ -diff(2),\ auto\ simp\ add:\ set\ -integrable\ -def\ set\ -integral\ -const[OF*]
algebra-simps intro!: integrable-mult-indicator)
         moreover have norm (\int x \in A. (f x - v) \partial N) \leq (\int x \in A. norm (f x))
(v) \partial N using * by (auto intro!: integral-norm-bound of N \lambda x. indicator A x
*_R (f x - v), THEN order-trans] integrable-mult-indicator integrable simp add:
set-lebesque-integral-def)
      ultimately have norm ((1 / measure N A) *_R set-lebesgue-integral N A f
-v) \leq set-lebesgue-integral N A (\lambda x. norm (f x - v)) / measure N A using asm
by (auto intro: divide-right-mono)
      also have ... < set-lebesgue-integral N A (\lambda x. r) / measure N A
        unfolding set-lebesque-integral-def
        using asm * integrable integrable-const AE-less measure-pos
     by (intro divide-strict-right-mono integral-less-AE[of - - A] integrable-mult-indicator)
         (fastforce simp add: dist-less dist-norm indicator-def)+
      also have \dots = r using * measure-pos by (simp add: set-integral-const)
      finally have dist ((1 / measure N A) *_R set-lebesgue-integral N A f) v < r
by (subst dist-norm)
    hence False using average [OF * measure-pos] by (metis\ ComplD\ dist-commute
in-mono mem-ball ball-in-Compl)
    thus ?thesis by fastforce
   qed
 note * = this
   fix b' assume b' \in B
   hence ball-subset-Compl: b' \subseteq -S and ball-radius-pos: \exists v \in D. \exists r > 0. b' =
ball v r using B-balls by (blast, fast)
 }
 \mathbf{note} ** = this
 hence emeasure N (f - b i \cap space N) = 0 for i by (cases b i = \{\}, simp)
(metis UnE singletonD * range-b[THEN eq-refl, THEN range-subsetD])
  hence emeasure N (\bigcup i. f - b i \cap space N) = 0 using open-b by (intro
emeasure-UN-eq-0) fastforce+
```

```
moreover have (\bigcup i. f - b \in Space N) = f - (\bigcup (range b)) \cap space N by
   ultimately have emeasure N (f - (-S) \cap space N) = 0 using Union-range-b
by argo
  hence AEx in N. fx \notin -S using open-Compl[OF assms(2)] by (intro AE-iff-measurable[THEN
iffD2], auto)
   thus ?case by force
qed (simp add: pred-sets2[OF borel-closed] assms(2))
lemma density-zero:
   fixes f::'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology, banach}\}
   assumes integrable M f
          and density-0: \bigwedge A. A \in sets \ M \Longrightarrow set-lebesgue-integral M \ A \ f = 0
   shows AE x in M. f x = 0
   using averaging-theorem[OF assms(1), of \{0\}] assms(2)
   by (simp add: scaleR-nonneq-nonneq)
This lemma shows that densities are unique in Banach spaces.
lemma density-unique-banach:
   fixes ff'::'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology, banach}\}
   assumes integrable M f integrable M f'
       and density-eq: \bigwedge A. A \in sets \ M \Longrightarrow set-lebesque-integral M A f = set-lebesque-integral
M A f'
   shows AE x in M. f x = f' x
proof-
   {
      fix A assume asm: A \in sets M
        hence LINT x|M. indicat-real A x *_R (f x - f' x) = 0 using density-eq
assms(1,2) by (simp\ add:\ set\ -lebesgue\ -integral\ -def\ algebra\ -simps\ Bochner\ -Integration\ .integral\ -diff\ |\ OF\ -lebesgue\ -integral\ -def\ algebra\ -simps\ Bochner\ -Integration\ .integral\ -diff\ |\ OF\ -lebesgue\ -integral\ -def\ algebra\ -simps\ Bochner\ -Integration\ .integral\ -diff\ |\ OF\ -lebesgue\ -integral\ -def\ algebra\ -simps\ Bochner\ -Integration\ .integral\ -diff\ |\ OF\ -lebesgue\ -lebesgue\
integrable-mult-indicator(1,1)])
   }
  thus ?thesis using density-zero[OF Bochner-Integration.integrable-diff[OF assms(1,2)]]
by (simp add: set-lebesgue-integral-def)
qed
lemma density-nonneg:
  fixes f::-\Rightarrow b::\{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, ordered\text{-}real\text{-}vector}\}
   assumes integrable M f
          and \bigwedge A. A \in sets \ M \Longrightarrow set-lebesgue-integral M \ A \ f \ge 0
      shows AE x in M. f x \ge 0
   using averaging-theorem [OF\ assms(1),\ of\ \{0..\},\ OF\ closed-atLeast]\ assms(2)
   by (simp add: scaleR-nonneg-nonneg)
corollary integral-nonneg-AE-eq-0-iff-AE:
    fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
   assumes f[measurable]: integrable M f and nonneg: AE x in M. 0 \le f x
   shows integral^L M f = 0 \longleftrightarrow (AE x in M. f x = 0)
proof
```

```
assume *: integral^L M f = 0
   fix A assume asm: A \in sets M
   have 0 \leq integral^L M (\lambda x. indicator A x *_R f x) using nonneg by (subst inte-
gral-zero[of M, symmetric], intro integral-mono-AE-banach integrable-mult-indicator
asm f integrable-zero, auto simp add: indicator-def)
  moreover have ... \leq integral^L M f using nonneg by (intro integral-mono-AE-banach
integrable-mult-indicator asm f, auto simp add: indicator-def)
  ultimately have set-lebesque-integral MAf = 0 unfolding set-lebesque-integral-def
\mathbf{using} * \mathbf{by} \; \mathit{force}
 }
 thus AE x in M. f x = 0 by (intro density-zero f, blast)
qed (auto simp add: integral-eq-zero-AE)
{\bf corollary}\ integral-eq-mono-AE-eq-AE:
  fixes f q :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
  assumes integrable M f integrable M q integral L M f = integral L M q AE x in
M. f x \leq g x
 shows AE x in M. f x = g x
proof -
  define h where h = (\lambda x. g x - f x)
  have AE \ x \ in \ M. \ h \ x = 0 unfolding h-def using assms by (subst inte-
gral-nonneg-AE-eq-0-iff-AE[symmetric]) auto
  then show ?thesis unfolding h-def by auto
qed
end
end
{\bf theory} \ {\it Conditional-Expectation-Banach}
imports\ HOL-Probability.\ Conditional-Expectation\ HOL-Probability.\ Independent-Family
Bochner-Integration-Supplement
begin
      Conditional Expectation in Banach Spaces
4
definition has-cond-exp:: 'a measure \Rightarrow 'a measure \Rightarrow ('a \Rightarrow 'b) \Rightarrow ('a \Rightarrow 'b::{real-normed-vector,
second-countable-topology\}) \Rightarrow bool where
  has-cond-exp M F f g = ((\forall A \in sets F. (\int x \in A. f x \partial M) = (\int x \in A. g x)
```

assumes $\bigwedge A$. $A \in sets \ F \Longrightarrow (\int x \in A. \ f \ x \ \partial M) = (\int x \in A. \ g \ x \ \partial M)$

 \land integrable M f \land integrable M g

 $\land g \in borel\text{-}measurable F$

 $\partial M))$

lemma has-cond-expI':

```
integrable M f
         integrable\ M\ g
         g \in borel-measurable F
 shows has\text{-}cond\text{-}exp\ M\ F\ f\ g
  using assms unfolding has-cond-exp-def by simp
lemma has-cond-expD:
 assumes has-cond-exp M F f g shows \bigwedge A. A \in sets \ F \Longrightarrow (\int x \in A. \ f \ x \ \partial M) = (\int x \in A. \ g \ x \ \partial M)
       integrable\ M\ f
       integrable M g
       g \in borel-measurable F
 using assms unfolding has-cond-exp-def by simp+
definition cond-exp :: 'a measure \Rightarrow 'a measure \Rightarrow ('a \Rightarrow 'b) \Rightarrow ('a \Rightarrow 'b::{banach,
second-countable-topology}) where
  cond-exp M F f = (if \exists g. has-cond-exp M F f g then (SOME g. has-cond-exp M
F f g) else (\lambda -. 0)
lemma borel-measurable-cond-exp[measurable]: cond-exp M F f \in borel-measurable
 by (metis cond-exp-def some I has-cond-exp-def borel-measurable-const)
lemma integrable-cond-exp[intro]: integrable M (cond-exp M F f)
 by (metis\ cond\text{-}exp\text{-}def\ has\text{-}cond\text{-}expD(3)\ integrable\text{-}zero\ some I)
{\bf lemma}\ set\text{-}integrable\text{-}cond\text{-}exp[intro]\text{:}
 assumes A \in sets M
 shows set-integrable M A (cond-exp M F f) using integrable-mult-indicator [OF
assms integrable-cond-exp, of F[f] by (auto simp add: set-integrable-def intro!: in-
tegrable-mult-indicator[OF\ assms\ integrable-cond-exp])
lemma has-cond-exp-self:
 assumes integrable M f
 shows has-cond-exp M (vimage-algebra (space M) f borel) ff
 using assms by (auto intro!: has-cond-expI' measurable-vimage-algebra1)
lemma has-cond-exp-sets-cong:
  assumes sets F = sets G
 shows has-cond-exp M F = has-cond-exp M G
 using assms unfolding has-cond-exp-def by force
lemma cond-exp-sets-cong:
 assumes sets F = sets G
 shows AE \ x \ in \ M. \ cond-exp \ M \ F \ f \ x = cond-exp \ M \ G \ f \ x
  by (intro AE-I2, simp add: cond-exp-def has-cond-exp-sets-cong[OF assms, of
M])
```

```
\begin{array}{l} \textbf{context} \ sigma-finite\text{-}subalgebra \\ \textbf{begin} \end{array}
```

 $\begin{tabular}{l} \textbf{lemma} borel-measurable-cond-exp'[measurable]: cond-exp M F $f \in borel-measurable M \\ M \end{tabular}$

by (metis cond-exp-def some I has-cond-exp-def borel-measurable-const subalg measurable-from-subalg)

```
lemma cond-exp-null:

assumes \nexists g. has-cond-exp M F f g

shows cond-exp M F f = (\lambda-. \theta)

unfolding cond-exp-def using assms by argo
```

We state the tower property of the conditional expectation in terms of the predicate has-cond-exp.

```
lemma has-cond-exp-nested-subalg:

fixes f:: 'a \Rightarrow 'b::\{second-countable-topology, banach\}

assumes subalgebra G F has-cond-exp M F f h has-cond-exp M G f h'

shows has-cond-exp M F h' h

by (intro has-cond-expI') (metis assms has-cond-expD in-mono subalgebra-def)+
```

The following lemma shows that the conditional expectation is unique as an element of L1, given that it exists.

```
lemma has-cond-exp-charact:
 fixes f :: 'a \Rightarrow 'b :: \{ second\text{-}countable\text{-}topology, banach \}
 assumes has-cond-exp M F f g
 shows has-cond-exp M F f (cond-exp M F f)
       AE \ x \ in \ M. \ cond\text{-}exp \ M \ F \ f \ x = g \ x
proof -
  show cond-exp: has-cond-exp M F f (cond-exp M F f) using assms some I
cond-exp-def by metis
 let ?MF = restr-to-subalg\ M\ F
 interpret sigma-finite-measure ?MF by (rule sigma-fin-subalg)
   fix A assume A \in sets ?MF
    then have [measurable]: A \in sets \ F using sets-restr-to-subalg[OF \ subalg] by
   have (\int x \in A. \ g \ x \ \partial ?MF) = (\int x \in A. \ g \ x \ \partial M) using assms subalg by (auto
simp add: integral-subalgebra2 set-lebesque-integral-def dest!: has-cond-expD)
    also have ... = (\int x \in A. \ cond\text{-}exp \ M \ F \ f \ x \ \partial M) using assms cond-exp by
(simp add: has-cond-exp-def)
   also have ... = (\int x \in A. \ cond\text{-}exp \ M \ F \ f \ x \ \partial ?MF) using subalg by (auto simp
add: integral-subalgebra2 set-lebesgue-integral-def)
   finally have (\int x \in A. \ g \ x \ \partial ?MF) = (\int x \in A. \ cond\text{-}exp \ M \ F \ f \ x \ \partial ?MF) by
simp
 hence AE x in ?MF. cond-exp M F f x = g x using cond-exp assms subalg by
(intro density-unique-banach, auto dest: has-cond-expD intro!: integrable-in-subalg)
```

```
then show AE \times in M. cond-exp M F f \times g \times using AE-restr-to-subalg[OF]
subalg by simp
qed
corollary cond-exp-charact:
 fixes f :: 'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology, banach\}
 assumes \bigwedge A. A \in sets \ F \Longrightarrow (\int x \in A. \ f \ x \ \partial M) = (\int x \in A. \ g \ x \ \partial M)
         integrable\ M\ f
         integrable M g
         g \in borel-measurable F
   shows AE x in M. cond-exp M F f x = g x
 by (intro has-cond-exp-charact has-cond-expI' assms) auto
corollary cond-exp-F-meas[intro, simp]:
 fixes f :: 'a \Rightarrow 'b :: \{ second\text{-}countable\text{-}topology, banach \}
 assumes integrable M f
        f \in borel-measurable F
   shows AE x in M. cond-exp M F f x = f x
 by (rule cond-exp-charact, auto intro: assms)
Congruence
lemma has-cond-exp-cong:
 assumes integrable M f \land x. x \in space M \Longrightarrow f x = g x has-cond-exp M F g h
 shows has\text{-}cond\text{-}exp\ M\ F\ f\ h
proof (intro\ has\text{-}cond\text{-}expI'[OF\text{-}assms(1)],\ goal\text{-}cases)
 case (1 A)
 hence set-lebesgue-integral MAf = set-lebesgue-integral MAg by (intro set-lebesgue-integral-cong)
(meson\ assms(2)\ subalg\ in-mono\ subalgebra-def\ sets.sets-into-space\ subalgebra-def
subsetD)+
 then show ?case using 1 assms(3) by (simp add: has-cond-exp-def)
qed (auto simp add: has-cond-expD[OF assms(3)])
lemma cond-exp-cong:
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach\}
 assumes integrable M f integrable M g \land x. x \in space M \Longrightarrow f x = g x
 shows AE x in M. cond-exp M F f x = cond-exp M F g x
proof (cases \exists h. has-cond-exp M F f h)
 case True
  then obtain h where h: has-cond-exp M F f h has-cond-exp M F g h using
has-cond-exp-cong assms by metis
 show ?thesis using h[THEN\ has\text{-}cond\text{-}exp\text{-}charact(2)] by fastforce
next
 {f case}\ {\it False}
 moreover have \nexists h. has-cond-exp M F q h using False has-cond-exp-cong assms
 ultimately show ?thesis unfolding cond-exp-def by auto
qed
lemma has-cond-exp-cong-AE:
```

```
assumes integrable M f AE x in M. f x = g x has-cond-exp M F g h
    shows has\text{-}cond\text{-}exp\ M\ F\ f\ h
   using assms(1,2) subalg subalgebra-def subset-iff
  by (intro has-cond-expI', subst set-lebesque-integral-cong-AE[OF - assms(1)]THEN
borel-measurable-integrable|\ borel-measurable-integrable(1)[OF\ has-cond-expD(2)]OF
assms(3)]])
       (fast\ intro:\ has-cond-expD[OF\ assms(3)]\ integrable-cong-AE-imp[OF\ -\ AE-symmetric])+
lemma has-cond-exp-cong-AE':
    assumes h \in borel-measurable F AE x in M. h x = h' x has-cond-exp M F f h'
    shows has\text{-}cond\text{-}exp\ M\ F\ f\ h
    using assms(1, 2) subalg subalgebra-def subset-iff
   using AE-restr-to-subalg2 [OF subalg assms(2)] measurable-from-subalg
  \textbf{by } (intro\ has\text{-}cond\text{-}expI'\ , subst\ set\text{-}lebesgue\text{-}integral\text{-}cong\text{-}AE[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}subalg}(1,1)[OF\text{-}measurable\text{-}from\text{-}s
subalq, OF - assms(1) has-cond-expD(4)[OF assms(3)]])
       (fast\ intro: has-cond-expD[OF\ assms(3)]\ integrable-conq-AE-imp[OF\ -\ -AE-symmetric])+
lemma cond-exp-cong-AE:
    fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
    assumes integrable M f integrable M g AE x in M. f x = g x
    shows AE x in M. cond-exp M F f x = cond-exp M F g x
proof (cases \exists h. has-cond-exp M F f h)
    case True
     then obtain h where h: has-cond-exp M F f h has-cond-exp M F q h using
has-cond-exp-cong-AE assms by (metis (mono-tags, lifting) eventually-mono)
    show ?thesis using h[THEN\ has\text{-}cond\text{-}exp\text{-}charact(2)] by fastforce
next
    case False
    moreover have \nexists h. has-cond-exp M F q h using False has-cond-exp-conq-AE
assms by auto
    ultimately show ?thesis unfolding cond-exp-def by auto
qed
lemma has-cond-exp-real:
   fixes f :: 'a \Rightarrow real
    assumes integrable M f
   \mathbf{shows}\ \mathit{has}\text{-}\mathit{cond}\text{-}\mathit{exp}\ \mathit{M}\ \mathit{F}\ \mathit{f}\ (\mathit{real}\text{-}\mathit{cond}\text{-}\mathit{exp}\ \mathit{M}\ \mathit{F}\ \mathit{f})
   by (intro has-cond-expI', auto intro!: real-cond-exp-intA assms)
lemma cond-exp-real[intro]:
    fixes f :: 'a \Rightarrow real
    assumes integrable M f
    shows AE \ x \ in \ M. \ cond-exp \ M \ F \ f \ x = real-cond-exp \ M \ F \ f \ x
    using has-cond-exp-charact has-cond-exp-real assms by blast
lemma cond-exp-cmult:
    fixes f :: 'a \Rightarrow real
    assumes integrable M f
    shows AE \ x \ in \ M. \ cond\text{-}exp \ M \ F \ (\lambda x. \ c * f \ x) \ x = c * cond\text{-}exp \ M \ F \ f \ x
```

using real-cond-exp-cmult[OF assms(1), of c] assms(1)[THEN cond-exp-real] assms(1)[THEN integrable-mult-right, THEN cond-exp-real, of c] by fastforce

4.1 Existence

Indicator functions

Addition

```
lemma has-cond-exp-indicator:
 assumes A \in sets \ M \ emeasure \ M \ A < \infty
  shows has-cond-exp M F (\lambda x. indicat-real A x *_R y) (\lambda x. real-cond-exp M F
(indicator\ A)\ x*_R\ y)
proof (intro has-cond-expI', goal-cases)
 case (1 B)
  have \int x \in B. (indicat-real A \times R y) \partial M = (\int x \in B. indicat-real A \times \partial M) *_R
y using assms by (intro set-integral-scaleR-left, meson 1 in-mono subalg subalge-
bra-def, blast
  also have ... = (\int x \in B. \ real\text{-}cond\text{-}exp \ M \ F \ (indicator \ A) \ x \ \partial M) *_R y \ using 1
assms by (subst real-cond-exp-intA, auto)
  also have ... = \int x \in B. (real-cond-exp M F (indicator A) x *_R y) \partial M using
assms by (intro set-integral-scaleR-left[symmetric], meson 1 in-mono subalg subal-
qebra-def, blast)
 finally show ?case.
next
 case 2
 then show ?case using integrable-scaleR-left integrable-real-indicator assms by
blast
\mathbf{next}
 case 3
 show ?case using assms by (intro integrable-scaleR-left, intro real-cond-exp-int,
blast+)
next
 case 4
 then show ?case by (intro borel-measurable-scaleR, intro Conditional-Expectation.borel-measurable-cond-exp.
\mathbf{qed}
lemma cond-exp-indicator[intro]:
 fixes y :: 'b:: \{second\text{-}countable\text{-}topology, banach\}
 assumes [measurable]: A \in sets \ M \ emeasure \ M \ A < \infty
 shows AE \times in M. cond-exp M F (\lambda x. indicat-real A \times *_R y) \times = cond-exp M F
(indicator\ A)\ x*_R\ y
proof -
 have AE \times in M. cond-exp M F (\lambda x. indicat-real A \times *_R y) \times = real-cond-exp M F
(indicator\ A)\ x*_R\ y\ \mathbf{using}\ has\text{-}cond\text{-}exp\text{-}indicator[OF\ assms]\ has\text{-}cond\text{-}exp\text{-}charact
by blast
 thus ?thesis using cond-exp-real[OF integrable-real-indicator, OF assms] by fast-
force
qed
```

```
lemma has-cond-exp-add:
  fixes fg :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
 assumes has-cond-exp M F f f' has-cond-exp M F g g'
 shows has-cond-exp M F (\lambda x. f x + g x) (\lambda x. f' x + g' x)
proof (intro has-cond-expI', goal-cases)
  case (1 A)
  have \int x \in A. (f x + g x) \partial M = (\int x \in A \cdot f x \partial M) + (\int x \in A \cdot g x \partial M) using
assms[THEN\ has-cond-expD(2)]\ subalg\ 1\  by (intro\ set-integral-add(2),\ auto\ simp
add: subalgebra-def set-integrable-def intro: integrable-mult-indicator)
  also have ... = (\int x \in A. \ f' \ x \ \partial M) + (\int x \in A. \ g' \ x \ \partial M) using assms[THEN]
has\text{-}cond\text{-}expD(1)[OF - 1]] by argo
 also have ... = \int x \in A. (f'x + g'x)\partial M using assms[THEN\ has-cond-expD(3)]
subalg 1 by (intro set-integral-add(2)[symmetric], auto simp add: subalgebra-def
set\text{-}integrable\text{-}def\ intro:\ integrable\text{-}mult\text{-}indicator)
 finally show ?case.
next
 case 2
 then show ?case by (metis Bochner-Integration.integrable-add assms has-cond-expD(2))
 then show ?case by (metis Bochner-Integration.integrable-add assms has-cond-expD(3))
next
  then show ?case using assms borel-measurable-add has-cond-expD(4) by blast
qed
lemma has-cond-exp-scaleR-right:
 fixes f :: 'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology,banach}\}
 assumes has-cond-exp M F f f'
 shows has-cond-exp M F (\lambda x. c *_R f x) (\lambda x. c *_R f' x)
 using has-cond-expD[OF assms] by (intro has-cond-expI', auto)
lemma cond-exp-scaleR-right:
 fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
 assumes integrable M f
 shows AE \ x \ in \ M. \ cond\text{-}exp \ M \ F \ (\lambda x. \ c *_R f \ x) \ x = c *_R \ cond\text{-}exp \ M \ F \ f \ x
proof (cases \exists f'. has-cond-exp M F f f')
  then show ?thesis using assms has-cond-exp-charact has-cond-exp-scaleR-right
by metis
\mathbf{next}
 case False
 show ?thesis
 proof (cases c = \theta)
   {\bf case}\  \, True
   then show ?thesis by simp
   case c-nonzero: False
   have \not\equiv f'. has-cond-exp M F (\lambda x. c *_R f x) f'
```

```
proof (standard, goal-cases)
     case 1
     then obtain f' where f': has-cond-exp M F (\lambda x. c *_R f x) f' by blast
      have has-cond-exp M F f (\lambda x. inverse c *_R f' x) using has-cond-expD[OF]
f'\ldot divideR-right[OF c-nonzero] assms by (intro has-cond-expI', auto)
     then show ?case using False by blast
   qed
   then show ?thesis using cond-exp-null[OF False] cond-exp-null by force
 qed
qed
lemma cond-exp-uminus:
 fixes f :: 'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology,banach\}
 assumes integrable M f
 shows AE x in M. cond-exp M F (\lambda x. - f x) x = - cond-exp M F f x
 using cond-exp-scaleR-right[OF assms, of -1] by force
corollary has-cond-exp-simple:
 fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
 assumes simple-function M f emeasure M \{y \in space M. f y \neq 0\} \neq \infty
 shows has-cond-exp M F f (cond-exp M F f)
 using assms
proof (induction rule: integrable-simple-function-induct)
  case (cong f g)
  then show ?case using has-cond-exp-cong by (metis (no-types, opaque-lifting)
Bochner-Integration.integrable-cong\ has-cond-exp(2)\ has-cond-exp-charact(1))
 case (indicator A y)
 then show ?case using has-cond-exp-charact[OF has-cond-exp-indicator] by fast
  case (add\ u\ v)
 then show ?case using has-cond-exp-add has-cond-exp-charact(1) by blast
qed
lemma cond-exp-contraction-real:
 fixes f :: 'a \Rightarrow real
 assumes integrable[measurable]: integrable M f
 shows AE \ x \ in \ M. \ norm \ (cond-exp \ M \ F \ f \ x) \le cond-exp \ M \ F \ (\lambda x. \ norm \ (f \ x)) \ x
proof-
 have int: integrable M (\lambda x. norm (f x)) using assms by blast
 have *: AE x in M. 0 \le cond\text{-}exp M F (\lambda x. norm (f x)) x using cond\text{-}exp\text{-}real[THEN]
AE-symmetric, OF integrable-norm [OF integrable] [real-cond-exp-ge-c [OF integrable-norm [OF]
integrable, of 0 norm-ge-zero by fastforce
  have **: A \in sets \ F \Longrightarrow \int x \in A. |f \ x| \ \partial M = \int x \in A. real-cond-exp M \ F \ (\lambda x).
norm (f x)) x \partial M for A unfolding real-norm-def using assms integrable-abs
real-cond-exp-intA by blast
 have norm-int: A \in sets \ F \Longrightarrow (\int x \in A. \ |f \ x| \ \partial M) = (\int x \in A. \ |f \ x| \ \partial M) for A
using assms by (intro nn-set-integral-eq-set-integral[symmetric], blast, fastforce)
```

```
(meson \ subalq \ subalqebra-def \ subset D)
```

have $AE \ x \ in \ M$. real-cond-exp $MF \ (\lambda x. \ norm \ (f \ x)) \ x \ge 0$ using int real-cond-exp-ge-c by force

hence cond-exp-norm-int: $A \in sets \ F \Longrightarrow (\int x \in A. \ real$ -cond- $exp \ M \ F \ (\lambda x. \ norm \ (f \ x)) \ x \ \partial M) = (\int {}^+x \in A. \ real$ -cond- $exp \ M \ F \ (\lambda x. \ norm \ (f \ x)) \ x \ \partial M)$ for A using assms by $(intro\ nn$ -set-integral-eq-set-integral[symmetric], blast, fastforce) $(meson\ subalg\ subalgebra$ - $def\ subset D)$

have $A \in sets \ F \Longrightarrow \int^+ x \in A$. $|f \ x| \partial M = \int^+ x \in A$. real-cond-exp $M \ F$ (λx . norm $(f \ x)$) $x \ \partial M$ for A using ** norm-int cond-exp-norm-int by (auto simp add: nn-integral-set-ennreal)

moreover have $(\lambda x. \ ennreal \ | f \ x |) \in borel-measurable \ M$ by measurable moreover have $(\lambda x. \ ennreal \ (real-cond-exp \ M \ F \ (\lambda x. \ norm \ (f \ x)) \ x)) \in borel-measurable \ F$ by measurable

ultimately have AE x in M. nn-cond-exp M F (λx . ennreal |f x|) x = real-cond-exp M F (λx . norm (f x)) x by ($intro\ nn$ -cond-exp-charact[$THEN\ AE$ -symmetric], auto)

hence AE x in M. nn-cond-exp M F (λx . ennreal |f x|) $x \leq cond$ -exp M F (λx . norm (f x)) x using cond-exp-real[OF int] by force

moreover have $AE \ x \ in \ M.$ $|real\text{-}cond\text{-}exp \ M \ F \ f \ x| = norm \ (cond\text{-}exp \ M \ F \ f \ x)$ unfolding real-norm-def using $cond\text{-}exp\text{-}real[OF \ assms] *$ by force

ultimately have $AE \ x \ in \ M$. $ennreal\ (norm\ (cond\text{-}exp\ M\ F\ f\ x)) \leq cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (f\ x))\ x\ using\ real\text{-}cond\text{-}exp\text{-}abs[OF\ assms[THEN\ borel\text{-}measurable\text{-}integrable]]}$ by fastforce

hence AE x in M. enn2real (ennreal (norm (cond-exp M F f $x))) <math>\leq$ enn2real (cond-exp M F (λx . norm (f x)) x) using ennreal-le-iff2 by force

thus ?thesis using * by fastforce qed

lemma cond-exp-contraction-simple:

fixes $f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}$

assumes simple-function M f emeasure M $\{y \in space M. f y \neq 0\} \neq \infty$

shows $AE \ x \ in \ M. \ norm \ (cond-exp \ M \ F \ f \ x) \le cond-exp \ M \ F \ (\lambda x. \ norm \ (f \ x)) \ x$ using assms

proof (induction rule: integrable-simple-function-induct)

case (cong f g)

hence ae: AE x in M. f x = q x by blast

hence AE x in M. cond-exp M F f x = cond-exp M F g x using cong has-cond-exp-simple by (subst cond-exp-cong-AE) (auto intro!: has-cond-expD(2))

hence $AE \ x \ in \ M. \ norm \ (cond-exp \ M \ F \ f \ x) = norm \ (cond-exp \ M \ F \ g \ x)$ by force

moreover have $AE \ x \ in \ M. \ cond-exp \ M \ F \ (\lambda x. \ norm \ (f \ x)) \ x = cond-exp \ M \ F \ (\lambda x. \ norm \ (g \ x)) \ x \ using \ ae \ cong \ has-cond-exp-simple \ by \ (subst \ cond-exp-cong-AE) \ (auto \ dest: \ has-cond-expD)$

ultimately show ?case using cong(6) by fastforce next

case (indicator A y)

hence $AE \ x$ in M. cond-exp $M \ F$ (λa . indicator $A \ a *_R y$) x = cond-exp $M \ F$

```
hence *: AE x in M. norm (cond-exp M F (\lambda a. indicat-real A a *_R y) x) \leq norm y
* cond-exp\ M\ F\ (\lambda x.\ norm\ (indicat-real\ A\ x))\ x\ {\bf using}\ cond-exp-contraction-real\ OF
integrable-real-indicator, OF indicator by fastforce
 have AE \times in M. norm y * cond-exp M F (\lambda x. norm (indicat-real A x)) x = norm
y * real\text{-}cond\text{-}exp \ M \ F \ (\lambda x. \ norm \ (indicat\text{-}real \ A \ x)) \ x \ \mathbf{using} \ cond\text{-}exp\text{-}real[OF]
integrable-real-indicator, OF indicator by fastforce
  moreover have AE x in M. cond-exp M F (\lambda x. norm y * norm (indicat-real
(A x)) x = real\text{-}cond\text{-}exp M F (\lambda x. norm y * norm (indicat\text{-}real A x)) x using 
indicator by (intro cond-exp-real, auto)
 ultimately have AE x in M. norm y * cond-exp M F (\lambda x. norm (indicat-real A))
x)) x = cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ y*norm\ (indicat\text{-}real\ A\ x))\ x\ using\ real\text{-}cond\text{-}exp\text{-}cmult[of
\lambda x. \ norm \ (indicat-real \ A \ x) \ norm \ y \ indicator \ \mathbf{by} \ fastforce
 moreover have (\lambda x. norm \ y * norm \ (indicat-real \ A \ x)) = (\lambda x. norm \ (indicat-real \ A \ x))
A x *_{B} y) by force
  ultimately show ?case using * by force
next
  case (add\ u\ v)
 have AE x in M. norm (cond-exp M F (\lambda a. u a + v a) x) = norm (cond-exp M
F \ u \ x + cond-exp M \ F \ v \ x) using has-cond-exp-charact(2)[OF has-cond-exp-add,
OF has-cond-exp-simple (1,1), OF add (1,2,3,4) by fastforce
  moreover have AE x in M. norm (cond-exp M F u x + cond-exp M F v x) \leq
norm (cond\text{-}exp \ M \ F \ u \ x) + norm (cond\text{-}exp \ M \ F \ v \ x) using norm\text{-}triangle\text{-}ineq
by blast
 moreover have AE \ x \ in \ M. \ norm \ (cond-exp \ M \ F \ u \ x) + norm \ (cond-exp \ M \ F \ v
x \le cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (u\ x))\ x + cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (v\ x))\ x using
add(6,7) by fastforce
 moreover have AE x in M. cond-exp M F (\lambda x. norm (u x)) x + cond-exp M F
(\lambda x. \ norm \ (v \ x)) \ x = cond\text{-}exp \ M \ F \ (\lambda x. \ norm \ (u \ x) + norm \ (v \ x)) \ x \ using \ in-
tegrable-simple-function [OF \ add(1,2)] integrable-simple-function [OF \ add(3,4)] by
(intro\ has\text{-}cond\text{-}exp\text{-}charact(2)[OF\ has\text{-}cond\text{-}exp\text{-}add[OF\ has\text{-}cond\text{-}exp\text{-}charact(1,1)],}
THEN AE-symmetric, auto intro: has-cond-exp-real)
 moreover have AE x in M. cond-exp MF (\lambda x. norm (u x) + norm (v x)) x =
cond-exp\ MF\ (\lambda x.\ norm\ (u\ x+v\ x))\ x\ using\ add(5)\ integrable-simple-function[OF]
add(1,2) integrable-simple-function [OF add(3,4)] by (intro cond-exp-cong, auto)
  ultimately show ?case by force
qed
lemma has-cond-exp-simple-lim:
   fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
  assumes integrable | measurable |: integrable | M f
      and \bigwedge i. simple-function M (s i)
      and \bigwedge i. emeasure M \{ y \in space M. \ s \ i \ y \neq 0 \} \neq \infty
      and \bigwedge x. \ x \in space \ M \Longrightarrow (\lambda i. \ s \ i \ x) \longrightarrow f \ x
      and \bigwedge x \ i. \ x \in space \ M \Longrightarrow norm \ (s \ i \ x) \le 2 * norm \ (f \ x)
  where strict-mono r has-cond-exp M F f (\lambda x. lim (\lambda i. cond-exp M F (s (r i))
x))
```

 $(indicator\ A)\ x*_R\ y\ \mathbf{by}\ blast$

```
AE \ x \ in \ M. \ convergent \ (\lambda i. \ cond-exp \ M \ F \ (s \ (r \ i)) \ x) proof -
```

have [measurable]: $(s\ i) \in borel$ -measurable M for i using assms(2) by $(simp\ add:\ borel$ -measurable-simple-function)

have integrable-s: integrable $M(\lambda x. s i x)$ for i using assms integrable-simple-function by blast

have integrable-4f: integrable M ($\lambda x.$ 4 * norm (fx)) using assms(1) by simp have integrable-2f: integrable M ($\lambda x.$ 2 * norm (fx)) using assms(1) by simp have integrable-2-cond-exp-norm-f: integrable M ($\lambda x.$ 2 * cond-exp M F ($\lambda x.$ norm (fx)) x) by fast

have emeasure M $\{y \in space \ M. \ s \ i \ y - s \ j \ y \neq 0\} \leq emeasure \ M \ \{y \in space \ M. \ s \ i \ y \neq 0\} \ for \ i \ j \ using <math>simple-functionD(2)[OF\ assms(2)]$ by $(intro\ order-trans[OF\ emeasure-mono\ emeasure-subadditive],\ auto)$

hence fin-sup: emeasure M { $y \in space M. \ s \ i \ y - s \ j \ y \neq 0$ } $\neq \infty$ **for** $i \ j \ using \ assms(3)$ **by** (metis (mono-tags) ennreal-add-eq-top linorder-not-less top.not-eq-extremum infinity-ennreal-def)

have emeasure M { $y \in space\ M$. norm ($s\ i\ y - s\ j\ y) \neq 0$ } \leq emeasure M { $y \in space\ M$. $s\ i\ y \neq 0$ } for $i\ j$ using $simple-functionD(2)[OF\ assms(2)]$ by (intro\ order-trans[OF\ emeasure-mono\ emeasure-subadditive], auto)

hence fin-sup-norm: emeasure M $\{y \in space M. norm (s i y - s j y) \neq 0\} \neq \infty$ for i j using assms(3) by (metis (mono-tags) ennreal-add-eq-top linorder-not-less top.not-eq-extremum infinity-ennreal-def)

have Cauchy: Cauchy $(\lambda n. \ s \ n \ x)$ if $x \in space \ M$ for x using assms(4) LIM-SEQ-imp-Cauchy that by blast

hence bounded-range-s: bounded (range $(\lambda n. \ s \ n \ x)$) if $x \in space \ M$ for x using that cauchy-imp-bounded by fast

have AE x in M. (λn . diameter $\{s \ i \ x \mid i. \ n \leq i\}$) $\longrightarrow 0$ using Cauchy cauchy-iff-diameter-tends-to-zero-and-bounded by fast

moreover have $(\lambda x.\ diameter\ \{s\ i\ x\ | i.\ n\leq i\})\in borel-measurable\ M\ for\ n$ using bounded-range-s borel-measurable-diameter by measurable

```
moreover have AE \ x \ in \ M. \ norm \ (diameter \ \{s \ i \ x \ | i. \ n \leq i\}) \leq 4 * norm \ (f \ x) \ for \ n proof —
```

fix x assume x: $x \in space M$

have diameter $\{s \ i \ x \ | i. \ n \leq i\} \leq 2 * norm \ (f \ x) + 2 * norm \ (f \ x)$ **by** (intro diameter-le, blast, subst dist-norm[symmetric], intro dist-triangle3[THEN order-trans, of 0], intro add-mono) (auto intro: $assms(5)[OF \ x]$)

hence norm (diameter $\{s \ i \ x \ | i.\ n \leq i\}$) $\leq 4 * norm (f \ x)$ using diameter-ge-0[OF bounded-subset[OF bounded-range-s], OF x, of $\{s \ i \ x \ | i.\ n \leq i\}$] by force

thus ?thesis by fast

```
qed
```

ultimately have diameter-tendsto-zero: ($\lambda n.\ LINT\ x|M.\ diameter\ \{s\ i\ x\mid i.\ n\leq i\}$) $\longrightarrow 0$ by (intro integral-dominated-convergence[OF borel-measurable-const[of\ 0] - integrable-4f, simplified]) (fast+)

have diameter-integrable: integrable M (λx . diameter $\{s \ i \ x \mid i. \ n \leq i\}$) for n using assms(1,5)

by (intro integrable-bound-diameter [OF bounded-range-s integrable-2f], auto)

have dist-integrable: integrable M (λx . dist (s i x) (s j x)) for i j using assms(5) dist-triangle 3 [of s i - 0, THEN order-trans, OF add-mono, of - 2 * norm (f -)] by (intro Bochner-Integration.integrable-bound [OF integrable-4f]) fastforce+

obtain N where *: LINT x|M. diameter $\{s \ i \ x \mid i. \ n \leq i\} < e \ \text{if} \ n \geq N \ \text{for}$ n using that order-tendsto-iff[THEN iffD1, OF diameter-tendsto-zero, unfolded eventually-sequentially] e-pos by presburger

fix i j x assume $asm: i \ge N j \ge N x \in space M$

have case-prod dist ' $(\{s\ i\ x\ | i.\ N \leq i\} \times \{s\ i\ x\ | i.\ N \leq i\}) = case-prod\ (\lambda i\ j.\ dist\ (s\ i\ x)\ (s\ j\ x))$ ' $(\{N...\} \times \{N...\})$ by fast

hence diameter $\{s \ i \ x \mid i.\ N \leq i\} = (SUP\ (i,j) \in \{N..\} \times \{N..\}.\ dist\ (s\ i \ x)\ (s\ j\ x))$ unfolding diameter-def by auto

moreover have $(SUP\ (i,j) \in \{N..\} \times \{N..\}.\ dist\ (s\ i\ x)\ (s\ j\ x)) \ge dist\ (s\ i\ x)\ (s\ j\ x)$ using asm bounded-imp-bdd-above [OF bounded-imp-dist-bounded, OF bounded-range-s] by (intro\ cSup-upper,\ auto)

ultimately have diameter $\{s \ i \ x \mid i. \ N \leq i\} \geq dist \ (s \ i \ x) \ (s \ j \ x)$ by presburger

hence LINT x|M. dist $(s\ i\ x)\ (s\ j\ x) < e\ \text{if}\ i \ge N\ j \ge N\ \text{for}\ i\ j\ \text{using}$ that $*\ \text{by}\ (intro\ integral-mono[OF\ dist-integrable\ diameter-integrable,\ THEN\ order.strict-trans1],\ blast+)$

moreover have LINT x|M. norm (cond-exp M F (s i) x – cond-exp M F (s j) x) \leq LINT x|M. dist (s i x) (s j x) for i j

proof -

have LINT x|M. norm (cond-exp M F (s i) x – cond-exp M F (s j) x) = LINT x|M. norm (cond-exp M F (s i) x + – 1 $*_R$ cond-exp M F (s j) x) unfolding dist-norm by simp

also have ... = LINT x|M. norm (cond-exp M F (λx . s i x - s j x) x) using has-cond-exp-charact(2)[OF has-cond-exp-add[OF - has-cond-exp-scaleR-right, OF has-cond-exp-charact(1,1), OF has-cond-exp-simple(1,1)[OF assms(2,3)]], THEN AE-symmetric, of i -1 j] by (intro integral-cong-AE) force+

also have $... \le LINT \ x | M. \ cond\text{-}exp \ M \ F \ (\lambda x. \ norm \ (s \ i \ x - s \ j \ x)) \ x \ using cond\text{-}exp\text{-}contraction\text{-}simple[OF - fin\text{-}sup, of i j] integrable\text{-}cond\text{-}exp \ assms(2) by (intro integral-mono-AE, fast+)$

also have ... = LINT x | M. norm (s i x - s j x) unfolding set-integral-space(1)[OF integrable-cond-exp, symmetric] set-integral-space[OF dist-integrable[unfolded dist-norm],

```
symmetric] \ \mathbf{by} \ (intro\ has\text{-}cond\text{-}expD(1)[OF\ has\text{-}cond\text{-}exp\text{-}simple[OF\text{-}fin\text{-}sup\text{-}norm]}, \\ symmetric]) \ (metis\ assms(2)\ simple\text{-}function\text{-}compose1\ simple\text{-}function\text{-}diff}, \\ metis\ sets.top\ subalg\ subalgebra\text{-}def)
```

finally show ?thesis unfolding dist-norm.

ged

ultimately show ?thesis using order.strict-trans1 by meson qed

then obtain r where strict-mono-r: strict-mono r and AE-Cauchy: AE x in M. Cauchy (λi . cond-exp M F (s (r i)) x) by ($rule\ cauchy$ -L1-AE-cauchy-subseq[OF integrable-cond-exp], auto)

hence ae-lim-cond-exp: $AE \ x \ in \ M$. $(\lambda n. \ cond-exp \ M \ F \ (s \ (r \ n)) \ x) \longrightarrow lim \ (\lambda n. \ cond-exp \ M \ F \ (s \ (r \ n)) \ x)$ using Cauchy-convergent-iff convergent-LIMSEQ-iff by fastforce

have cond-exp-bounded: AE x in M. norm (cond-exp M F (s (r n)) x) \leq cond-exp M F (λx . 2 * norm (f x)) x **for** n

proof -

have $AE \ x \ in \ M. \ norm \ (cond\text{-}exp \ M \ F \ (s \ (r \ n)) \ x) \leq cond\text{-}exp \ M \ F \ (\lambda x. \ norm \ (s \ (r \ n) \ x)) \ x \ \mathbf{by} \ (rule \ cond\text{-}exp\text{-}contraction\text{-}simple[OF \ assms(2,3)])$

moreover have AE x in M. real-cond-exp M F $(\lambda x. norm (s (r n) x))$ $x \le real-cond-exp$ M F $(\lambda x. 2 * norm (f x))$ x **using** integrable-s integrable-s integrable-s assms(5) by (intro real-cond-exp-mono, auto)

ultimately show ?thesis using cond-exp-real[OF integrable-norm, OF integrable-s, of r n] cond-exp-real[OF integrable-2f] by force qed

have lim-integrable: integrable M (λx . lim (λi . cond-exp M F (s (r i)) x)) by (intro integrable-dominated-convergence[OF - borel-measurable-cond-exp' integrable-cond-exp ae-lim-cond-exp cond-exp-bounded], simp)

fix A assume A-in-sets-F: $A \in sets F$

have $AE \ x \ in \ M. \ norm \ (indicator \ A \ x *_R \ cond-exp \ M \ F \ (s \ (r \ n)) \ x) \leq cond-exp \ M \ F \ (\lambda x. \ 2 * norm \ (f \ x)) \ x \ {\bf for} \ n$ proof -

have $AE \ x \ in \ M. \ norm \ (indicator \ A \ x *_R \ cond-exp \ M \ F \ (s \ (r \ n)) \ x) \leq norm \ (cond-exp \ M \ F \ (s \ (r \ n)) \ x) \ unfolding \ indicator-def \ by \ simp$

thus ?thesis using cond-exp-bounded[of n] by force

hence lim-cond-exp-int: $(\lambda n. \ LINT \ x:A|M. \ cond-exp \ M \ F \ (s \ (r \ n)) \ x) \longrightarrow LINT \ x:A|M. \ lim \ (\lambda n. \ cond-exp \ M \ F \ (s \ (r \ n)) \ x)$

 $\begin{tabular}{l} \textbf{using} \ ae-lim-cond-exp \ measurable-from-subalg[OF \ subalg \ borel-measurable-indicator, \ OF \ A-in-sets-F] \ cond-exp-bounded \end{tabular}$

unfolding set-lebesgue-integral-def

 $\mathbf{by} \ (intro \ integral-dominated-convergence [OF \ borel-measurable-scale R \ borel-measurable-scale R \ integrable-cond-exp]) \ (fastforce \ simp \ add: \ tendsto-scale R)+$

have $AE\ x\ in\ M.\ norm\ (indicator\ A\ x*_R\ s\ (r\ n)\ x) \leq 2*\ norm\ (f\ x)\ {\bf for}\ n$ proof -

have AE x in M. norm (indicator A $x *_R s (r n) x) \leq norm (s (r n) x)$

```
unfolding indicator-def by simp
      thus ?thesis using assms(5)[of - r n] by fastforce
   qed
   hence lim-s-int: (\lambda n. \ LINT \ x:A|M. \ s \ (r \ n) \ x) \longrightarrow LINT \ x:A|M. \ f \ x
    using measurable-from-subalq[OF subalq borel-measurable-indicator, OF A-in-sets-F]
LIMSEQ-subseq-LIMSEQ[OF\ assms(4)\ strict-mono-r]\ assms(5)
      unfolding set-lebesque-integral-def comp-def
    \textbf{by } (intro\ integral-dominated-convergence [OF\ borel-measurable-scaleR\ borel-measurable-scaleR]
integrable-2f]) (fastforce\ simp\ add:\ tendsto-scaleR)+
     have LINT x:A|M. lim (\lambda n. cond\text{-}exp \ M \ F \ (s \ (r \ n)) \ x) = lim \ (\lambda n. \ LINT
x:A|M.\ cond-exp\ M\ F\ (s\ (r\ n))\ x) using limI[OF\ lim-cond-exp-int] by argo
   also have ... = lim (\lambda n. LINT x: A|M. s (r n) x) using has\text{-}cond\text{-}expD(1)[OF]
has\text{-}cond\text{-}exp\text{-}simple[OF\ assms(2,3)]\ A\text{-}in\text{-}sets\text{-}F,\ symmetric}] by presburger
   also have ... = LINT x:A|M. fx using limI[OF lim-s-int] by argo
   finally have LINT x:A|M. lim (\lambda n. cond\text{-}exp \ M \ F \ (s \ (r \ n)) \ x) = LINT \ x:A|M.
fx.
  hence has-cond-exp M F f (\lambda x. lim (\lambda i. cond-exp M F (s (r i)) x)) using
assms(1) lim-integrable by (intro has-cond-expI', auto)
  thus thesis using AE-Cauchy Cauchy-convergent strict-mono-r by (auto intro!:
that)
\mathbf{qed}
corollary has-cond-expI:
  fixes f :: 'a \Rightarrow 'b :: \{ second\text{-}countable\text{-}topology, banach \}
  assumes integrable M f
  shows has-cond-exp M F f (cond-exp M F f)
proof -
 obtain s where s-is: \bigwedge i. simple-function M (s i) \bigwedge i. emeasure M {y \in space M.
\{s \mid y \neq 0\} \neq \infty \land x. \ x \in space \ M \Longrightarrow (\lambda i. \ s \mid x) \longrightarrow fx \land x \ i. \ x \in space \ M \Longrightarrow \{\lambda i. \ s \mid x\} 
norm\ (s\ i\ x) \le 2*norm\ (f\ x) using integrable-implies-simple-function-sequence [OF]
assms] by blast
 show ?thesis using has-cond-exp-simple-lim[OF assms s-is] has-cond-exp-charact(1)
by metis
qed
4.2
        Properties
lemma cond-exp-nested-subalg:
  fixes f :: 'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology,banach}\}
  assumes integrable M f subalgebra M G subalgebra G F
 shows AE \xi in M. cond-exp M F f \xi = cond-exp M F (cond-exp M G f) \xi
 {f using}\ has\text{-}cond\text{-}expI\ assms\ sigma\text{-}finite\text{-}subalgebra\text{-}def\ {f by}\ (auto\ intro!:\ has\text{-}cond\text{-}exp\text{-}nested\text{-}subalg}\ [THEN]
has\text{-}cond\text{-}exp\text{-}charact(2), THEN\ AE\text{-}symmetric]\ sigma\text{-}finite\text{-}subalgebra.has\text{-}cond\text{-}expI[OF]
sigma-finite-subalgebra.intro[OF\ assms(2)]]\ nested-subalg-is-sigma-finite)
lemma cond-exp-set-integral:
  fixes f :: 'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology,banach}\}
```

```
shows (\int x \in A. fx \partial M) = (\int x \in A. cond\text{-}exp M F fx \partial M)
  using has\text{-}cond\text{-}expD(1)[OF\ has\text{-}cond\text{-}expI,\ OF\ assms] by argo
lemma cond-exp-add:
  fixes f :: 'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology,banach}\}
  assumes integrable M f integrable M g
  shows AE x in M. cond-exp M F (\lambda x. f x + g x) x = cond-exp M F f x +
cond-exp M F g x
 using has-cond-exp-add OF has-cond-expI(1,1), OF assms, THEN has-cond-exp-charact (2)
lemma cond-exp-diff:
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
  assumes integrable M f integrable M q
  shows AE x in M. cond-exp M F (\lambda x. f x - g x) x = cond-exp M F f x -
cond-exp M F g x
 using has\text{-}cond\text{-}exp\text{-}add[OF\text{-}has\text{-}cond\text{-}exp\text{-}scaleR\text{-}right, OF }has\text{-}cond\text{-}exp}I(1,1),
OF assms, THEN has-cond-exp-charact(2), of -1] by simp
lemma cond-exp-diff':
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
  assumes integrable M f integrable M g
  shows AE \ x \ in \ M. \ cond-exp \ M \ F \ (f - g) \ x = cond-exp \ M \ F \ f \ x - cond-exp \ M
  unfolding fun-diff-def using assms by (rule cond-exp-diff)
lemma cond-exp-scaleR-left:
  fixes f :: 'a \Rightarrow real
  assumes integrable M f
  shows AE x in M. cond-exp M F (\lambda x. f x *_R c) x = cond-exp M F f x *_R c
  using cond-exp-set-integral [OF assms] subalg assms unfolding subalgebra-def
  by (intro cond-exp-charact,
     subst set-integral-scaleR-left, blast, intro assms,
     subst set-integral-scaleR-left, blast, intro integrable-cond-exp)
     auto
lemma cond-exp-contraction:
  fixes f::'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology, banach\}
  assumes integrable M f
 shows AE \ x \ in \ M. \ norm \ (cond-exp \ M \ F \ f \ x) \le cond-exp \ M \ F \ (\lambda x. \ norm \ (f \ x))
proof -
 obtain s where s: \bigwedge i. simple-function M (s i) \bigwedge i. emeasure M \{y \in space M.
\{s \ i \ y \neq 0\} \neq \infty \ \land x. \ x \in space \ M \Longrightarrow (\lambda i. \ s \ i \ x) \longrightarrow f \ x \ \land i \ x. \ x \in space \ M
```

assumes integrable $M f A \in sets F$

```
\implies norm (s \ i \ x) \le 2 * norm (f \ x)
by (blast intro: integrable-implies-simple-function-sequence [OF assms])
```

obtain r where r: strict-mono r and has-cond-exp M F f $(\lambda x. lim$ $(\lambda i. cond$ -exp M F (s(ri)) x)) AE x in M. $(\lambda i. cond$ -exp M F (s(ri)) x) $\longrightarrow lim$ $(\lambda i. cond$ -exp M F (s(ri)) x)

using has-cond-exp-simple-lim[OF assms s] unfolding convergent-LIMSEQ-iff by blast

hence r-tendsto: AE x in M. (λi . cond-exp M F (s (r i)) x) \longrightarrow cond-exp M F f x using has-cond-exp-charact(2) by force

have norm-s-r: $\land i$. simple-function M (λx . norm (s (r i) x)) $\land i$. emeasure M { $y \in space\ M$. norm (s (r i) y) $\neq 0$ } $\neq \infty \land x$. $x \in space\ M \Longrightarrow (\lambda i.\ norm\ (<math>s$ (r i) x)) \longrightarrow norm (f x) $\land i$ x. $x \in space\ M \Longrightarrow$ norm ($norm\ (<math>s$ (r i) x)) $\leq 2 *$ norm ($norm\ (fx)$)

using s **by** ($auto\ intro:\ LIMSEQ$ -subseq- $LIMSEQ[OF\ tendsto-norm\ r,\ unfolded\ comp-def]\ simple-function-compose1)$

obtain r' where r': strict-mono r' and has-cond-exp M F $(\lambda x. norm (f x)) (\lambda x. lim (\lambda i. cond-exp <math>M$ F $(\lambda x. norm (s (r (r' i)) x)) x)) AE x in <math>M$. ($\lambda i. cond$ -exp M F $(\lambda x. norm (s (r (r' i)) x)) x) <math>\longrightarrow lim (\lambda i. cond$ -exp M F $(\lambda x. norm (s (r (r' i)) x)) x)$ using has-cond-exp-simple-lim[OF integrable-norm norm-s-r, OF assms] unfolding convergent-LIMSEQ-iff by blast

hence r'-tendsto: $AE \ x \ in \ M. \ (\lambda i. \ cond\text{-}exp \ M \ F \ (\lambda x. \ norm \ (s \ (r \ (r' \ i)) \ x)) \ x)$ $\longrightarrow cond\text{-}exp \ M \ F \ (\lambda x. \ norm \ (f \ x)) \ x \ using \ has-cond\text{-}exp-charact(2) \ by \ force$

have $AE \ x \ in \ M$. $\forall i. \ norm \ (cond\text{-}exp \ M \ F \ (s \ (r \ (r' \ i))) \ x) \leq cond\text{-}exp \ M \ F \ (\lambda x. \ norm \ (s \ (r \ (r' \ i)) \ x)) \ x \ using \ s \ by \ (auto \ intro: \ cond\text{-}exp\text{-}contraction\text{-}simple \ simp \ add: } AE\text{-}all\text{-}countable})$

moreover have $AE \ x \ in \ M. \ (\lambda i. \ norm \ (cond\text{-}exp \ M \ F \ (s \ (r \ (r' \ i))) \ x)) \longrightarrow norm \ (cond\text{-}exp \ M \ F \ f \ x) \ using \ r\text{-}tendsto \ LIMSEQ\text{-}subseq\text{-}LIMSEQ[OF \ tend-sto-norm \ r', unfolded \ comp\text{-}def]}$ by fast

ultimately show ?thesis using LIMSEQ-le r'-tendsto by fast qed

 $\mathbf{lemma}\ cond\text{-}exp\text{-}measurable\text{-}mult$:

```
fixes f g :: 'a \Rightarrow real
```

assumes [measurable]: integrable M ($\lambda x. fx * gx$) integrable M $gf \in borel$ -measurable F

shows integrable M ($\lambda x. f x * cond\text{-}exp M F g x$)

 $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ (\lambda x.\ f\ x\ *\ g\ x)\ x = f\ x\ *\ cond\text{-}exp\ M\ F\ g\ x$ $\mathbf{proof}-$

show integrable: integrable M ($\lambda x. fx * cond-exp\ MFgx$) using $cond-exp-real[OF\ assms(2)]$ by (intro integrable-cong-AE-imp[OF real-cond-exp-intg(1), OF assms(1,3) $assms(2)[THEN\ borel-measurable-integrable]]$ measurable-from-subalg[OF\ subalg]) auto

interpret sigma-finite-measure restr-to-subalg M F by (rule sigma-fin-subalg)

```
fix A assume asm: A \in sets F
   hence asm': A \in sets \ M using subalg by (fastforce \ simp \ add: \ subalgebra-def)
  have set-lebesque-integral M A (cond-exp M F (\lambda x. f x * g x)) = set-lebesque-integral
M A (\lambda x. f x * g x) by (simp add: cond-exp-set-integral [OF assms(1) asm])
     also have ... = set-lebesque-integral M A (\lambda x. f x * real-cond-exp M F g
x) using borel-measurable-times [OF borel-measurable-indicator [OF asm] assms(3)]
borel-measurable-integrable[OF\ assms(2)]\ integrable-mult-indicator[OF\ asm'\ assms(1)]
by (fastforce simp add: set-lebesque-integral-def mult. assoc[symmetric] intro: real-cond-exp-intq(2)[symmetric])
    also have ... = set-lebesgue-integral M A (\lambda x. f x * cond\text{-}exp M F g x) using
cond-exp-real[OF\ assms(2)]\ asm'\ borel-measurable-cond-exp'\ borel-measurable-cond-exp2
measurable-from-subalg [OF subalg assms(3)] by (auto simp add: set-lebesgue-integral-def
intro: integral-cong-AE)
   finally have set-lebesgue-integral M A (cond-exp M F (\lambda x. f x * g x)) = \int x \in A.
(f x * cond\text{-}exp M F g x)\partial M.
  }
  hence AE x in restr-to-subalg M F. cond-exp M F (\lambda x. f x * g x) x = f
x * cond-exp M F g x  by (intro density-unique-banach integrable-cond-exp inte-
grable integrable-in-subalg subalg, measurable, simp add: set-lebesgue-integral-def
integral-subalgebra2[OF\ subalg]\ sets-restr-to-subalg[OF\ subalg])
  thus AE x in M. cond-exp M F (\lambda x. f x * g x) x = f x * cond-exp M F g x by
(rule\ AE-restr-to-subalg[OF\ subalg])
qed
\mathbf{lemma}\ cond\text{-}exp\text{-}measurable\text{-}scaleR:
 fixes f :: 'a \Rightarrow real and g :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
 assumes [measurable]: integrable M (\lambda x. fx *_R qx) integrable M q f \in borel-measurable
F
 shows integrable M (\lambda x. f x *_R cond\text{-}exp M F g x)
       AE \ x \ in \ M. \ cond-exp \ M \ F \ (\lambda x. \ f \ x *_R \ g \ x) \ x = f \ x *_R \ cond-exp \ M \ F \ g \ x
proof -
 let ?F = restr-to-subalq M F
 have subalg': subalgebra M (restr-to-subalg M F) by (metis sets-eq-imp-space-eq
sets-restr-to-subalg subalg subalgebra-def)
  fix z assume asm[measurable]: integrable\ M\ (\lambda x.\ z\ x*_R\ g\ x)\ z\in borel-measurable
?F
    hence asm'[measurable]: z \in borel-measurable F using measurable-in-subalg'
subalg by blast
    have integrable M (\lambda x. z x *_R cond\text{-}exp M F g x) LINT x|M. z x *_R g x =
LINT \ x|M. \ z \ x *_R \ cond-exp \ M \ F \ g \ x
   proof -
     obtain s where s-is: \bigwedge i. simple-function ?F(s i) \bigwedge x. x \in space ?F \Longrightarrow (\lambda i.
          \longrightarrow z \ x \ \bigwedge i \ x. \ x \in space \ ?F \Longrightarrow norm \ (s \ i \ x) \le 2 * norm \ (z \ x) \ \mathbf{using}
```

have s-scaleR-g-tendsto: AE x in M. $(\lambda i. \ s \ i \ x *_R g \ x) \longrightarrow z \ x *_R g \ x$

borel-measurable-implies-sequence-metric [OF asm(2), of 0] by force

```
using s-is(2) by (simp add: space-restr-to-subalg tendsto-scaleR)
have s-scaleR-cond-exp-g-tendsto: AE x in ?F. (\lambda i. s i x *<sub>R</sub> cond-exp M F g
x) \longrightarrow z x *<sub>R</sub> cond-exp M F g x using s-is(2) by (simp add: tendsto-scaleR)
```

have s-scaleR-g-meas: $(\lambda x.\ s\ i\ x*_R\ g\ x)\in borel$ -measurable M for i using s-is(1)[THEN borel-measurable-simple-function, THEN subalg'[THEN measurable-from-subalg]] by simp

have s-scaleR-cond-exp-g-meas: $(\lambda x.\ s\ i\ x*_R\ cond-exp\ M\ F\ g\ x)\in borel$ -measurable ?F for i using s- $is(1)[THEN\ borel$ -measurable-simple-function] measurable-in-subalg[OF subalg borel-measurable-cond-exp] by (fastforce intro: borel-measurable-scaleR)

```
have s-scaleR-g-AE-bdd: AE x in M. norm (s i x *_R g x) \leq 2 * norm (z x *_R g x) for i using s-is(3) by (fastforce simp add: space-restr-to-subalg mult.assoc[symmetric] mult-right-mono) { fix i
```

have asm: integrable M (λx . norm (z x) * norm (g x)) using asm(1)[THEN integrable-norm] by simp

have $\overrightarrow{AE} x$ in ?F. norm $(s \ i \ x *_R \ cond\text{-}exp \ M \ F \ g \ x) \le 2 * norm \ (z \ x) * norm \ (cond\text{-}exp \ M \ F \ g \ x)$ using s-is(3) by $(fastforce \ simp \ add: \ mult-mono)$

moreover have $AE \ x \ in \ ?F. \ norm \ (z \ x) * cond-exp \ MF \ (\lambda x. \ norm \ (g \ x)) \ x = cond-exp \ MF \ (\lambda x. \ norm \ (z \ x) * norm \ (g \ x)) \ x \ \mathbf{by} \ (rule \ cond-exp-measurable-mult(2)[THEN \ AE-symmetric, \ OF \ asm \ integrable-norm, \ OF \ assms(2), \ THEN \ AE-restr-to-subalg2[OF \ subalg]], \ auto)$

ultimately have $AE \ x \ in \ ?F. \ norm \ (s \ i \ x *_R \ cond-exp \ M \ F \ g \ x) \le 2 * cond-exp \ M \ F \ (\lambda x. \ norm \ (z \ x *_R \ g \ x)) \ x \ using \ cond-exp-contraction[OF \ assms(2), THEN \ AE-restr-to-subalg2[OF \ subalg]] \ order-trans[OF - mult-mono] \ by \ fastforce$ } note s-scaleR-cond-exp-g-AE-bdd = this

```
{
fix i
```

have s-meas-M[measurable]: $s \ i \in borel$ -measurable M by (meson borel-measurable-simple-function measurable-from-subalq s-is(1) subalq')

have s-meas-F[measurable]: $s \ i \in borel$ -measurable F by $(meson\ borel$ -measurable-simple-function measurable-in-subalg' s-is $(1)\ subalg)$

have s-scaleR-eq: s i $x *_R h$ $x = (\sum y \in s$ i 'space M. (indicator (s i - ' $\{y\} \cap space M$) $x *_R y$) $*_R h$ x) if $x \in space M$ for x and h :: ' $a \Rightarrow$ 'b using simple-function-indicator-representation [OF s-is(1), of x i] that unfolding space-restr-to-subalg scaleR-left.sum[of - - h x, symmetric] by presburger

have LINT x|M. s i $x *_R g$ x = LINT x|M. $(\sum y \in s$ i 'space M. indicator (s i - ' $\{y\}$ \cap space M) $x *_R y *_R g$ x) using s-scaleR-eq by (intro Bochner-Integration.integral-cong) auto

also have ... = $(\sum y \in s \ i \ `space M. \ LINT \ x|M. \ indicator \ (s \ i - `\{y\} \cap space M) \ x *_R y *_R g \ x)$ by $(intro \ Bochner-Integration.integral-sum \ in-space M)$

```
tegrable-mult-indicator[OF - integrable-scaleR-right] assms(2)) simp
       also have ... = (\sum y \in s \ i \ `space M. \ y *_R set-lebesgue-integral M \ (s \ i - `
\{y\} \cap space\ M)\ g)\ \mathbf{by}\ (simp\ only:\ set-lebesgue-integral-def[symmetric])\ simp
      also have ... = (\sum y \in s \ i \ 'space \ M. \ y *_R set-lebesgue-integral \ M \ (s \ i - ' \{y\}
\cap space M) (cond-exp M F g)) using assms(2) subalg borel-measurable-vimage[OF]
s-meas-F] by (subst cond-exp-set-integral, auto simp add: subalgebra-def)
     also have ... = (\sum y \in s \ i \ `space \ M. \ LINT \ x | M. \ indicator \ (s \ i - `\{y\} \cap space
M) \ x *_R y *_R cond-exp M F g x)  by (simp only: set-lebesgue-integral-def[symmetric])
simp
     also have ... = LINT x|M. (\sum y \in s \ i 'space M. indicator (s \ i - '\{y\} \cap space
M) x *_R y *_R cond-exp M F g x by (intro Bochner-Integration.integral-sum[symmetric]
integrable-mult-indicator[OF - integrable-scaleR-right]) auto
       also have ... = LINT x|M. s i x *_R cond-exp M F g x using s-scaleR-eq
by (intro Bochner-Integration.integral-cong) auto
      finally have LINT x|M. s i x *_R g x = LINT x|?F. s i x *_R cond-exp M F
g \times y = 1 by (simp \ add: integral-subalgebra2[OF \ subalg])
     note integral-s-eq = this
   show integrable M (\lambda x. zx *_R cond-exp M F g x) using s-scaleR-cond-exp-g-meas
asm(2) borel-measurable-cond-exp' by (intro integrable-from-subalg[OF subalg] in-
tegrable-cond-exp integrable-dominated-convergence [OF - - - s-scaleR-cond-exp-q-tendsto
s-scaleR-cond-exp-g-AE-bdd]) (auto intro: measurable-from-subalg[OF subalg] inte-
grable-in-subalg measurable-in-subalg subalg)
      have (\lambda i. \ LINT \ x | M. \ s \ i \ x *_R \ g \ x) \longrightarrow LINT \ x | M. \ z \ x *_R \ g \ x \ using
s-scaleR-g-meas asm(1)[THEN integrable-norm] asm' borel-measurable-cond-exp'
by (intro integral-dominated-convergence OF - - - s-scaleR-g-tends to s-scaleR-g-AE-bdd))
(auto intro: measurable-from-subalg[OF subalg])
       moreover have (\lambda i. \ LINT \ x|?F. \ s \ i \ x *_R \ cond-exp \ M \ F \ g \ x) —
LINT x \mid ?F. z \mid x \mid *_R cond-exp \mid M \mid F \mid g \mid x  using s-scaleR-cond-exp-g-meas asm(2)
borel-measurable-cond-exp' by (intro integral-dominated-convergence OF - - - s-scaleR-cond-exp-g-tendsto
s-scaleR-cond-exp-g-AE-bdd]) (auto intro: measurable-from-subalg[OF subalg] inte-
grable-in-subalq measurable-in-subalq subalq)
      ultimately show LINT x|M. z x *_R g x = LINT x|M. z x *_R cond-exp
M F g x using integral-s-eq using subalg by (simp add: LIMSEQ-unique inte-
gral-subalgebra2)
   qed
 note * = this
  show integrable M (\lambda x. f x *_R cond\text{-}exp M F g x) using * assms measur-
able-in-subalg[OF subalg] by blast
  {
```

```
fix A assume asm: A \in F
    hence integrable M (\lambda x. indicat-real A x *_R f x *_R g x) using subalg by
(fastforce\ simp\ add:\ subalgebra-def\ intro!:\ integrable-mult-indicator\ assms(1))
    hence set-lebesque-integral M A (\lambda x. f x *_R g x) = set-lebesque-integral M A
(\lambda x. f x *_R cond\text{-}exp M F q x) unfolding set-lebesque-integral-def using asm by
(auto intro!: * measurable-in-subalg[OF subalg])
  }
  thus AE x in M. cond-exp M F (\lambda x. f x *<sub>R</sub> q x) x = f x *<sub>R</sub> cond-exp M F q x
using borel-measurable-cond-exp by (intro cond-exp-charact, auto intro!: * assms
measurable-in-subalg[OF\ subalg])
qed
lemma cond-exp-sum [intro, simp]:
  fixes f :: 't \Rightarrow 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
 assumes [measurable]: \bigwedge i. integrable M (f i)
 shows AE \ x \ in \ M. \ cond\text{-}exp \ M \ F \ (\lambda x. \ \sum i \in I. \ f \ i \ x) \ x = (\sum i \in I. \ cond\text{-}exp \ M \ F
(f i) x
proof (rule has-cond-exp-charact, intro has-cond-expI')
  fix A assume [measurable]: A \in sets F
  then have A-meas [measurable]: A \in sets \ M by (meson subsetD subalg subalge-
bra-def)
  have (\int x \in A. (\sum i \in I. f i x) \partial M) = (\int x. (\sum i \in I. indicator A x *_R f i x) \partial M)
unfolding set-lebesgue-integral-def by (simp add: scaleR-sum-right)
 also have ... = (\sum i \in I. (\int x. indicator A x *_R f i x \partial M)) using assms by (auto
intro!: Bochner-Integration.integral-sum integrable-mult-indicator)
 also have ... = (\sum i \in I. (\int x. indicator A x *_R cond-exp M F (f i) x \partial M)) using
cond-exp-set-integral[OF assms] by (simp add: set-lebesgue-integral-def)
  also have ... = (\int x. (\sum i \in I. indicator \ A \ x *_R cond-exp \ M \ F \ (f \ i) \ x)\partial M)
using assms by (auto intro!: Bochner-Integration.integral-sum[symmetric] inte-
grable-mult-indicator)
 also have ... = (\int x \in A. (\sum i \in I. cond\text{-}exp\ M\ F\ (f\ i)\ x) \partial M) unfolding set-lebesgue-integral-def
by (simp add: scaleR-sum-right)
 finally show (\int x \in A. (\sum i \in I. fi x) \partial M) = (\int x \in A. (\sum i \in I. cond\text{-}exp M F (fi))
x)\partial M) by auto
qed (auto simp add: assms integrable-cond-exp)
```

4.3 Linearly Ordered Banach Spaces

In this subsection we show monotonicity results concerning the conditional expectation operator.

```
lemma cond-exp-gr-c:
    fixes f:: 'a \Rightarrow 'b:: \{second-countable-topology, banach, linorder-topology, or-
dered-real-vector\}
    assumes integrable\ M\ f\ AE\ x\ in\ M.\ f\ x>c
    shows AE\ x\ in\ M.\ cond-exp\ M\ F\ f\ x>c

proof -
    define X where X=\{x\in space\ M.\ cond-exp\ M\ F\ f\ x\leq c\}
    have [measurable]:\ X\in sets\ F unfolding X-def by measurable\ (metis\ sets.top)
```

```
subalq\ subalqebra-def)
  hence X-in-M: X \in sets \ M using sets-restr-to-subalg subalgebra-def by
 have emeasure M X = 0
  proof (rule ccontr)
   assume emeasure M X \neq 0
   have emeasure (restr-to-subalg M F) X = emeasure M X by (simp add: emea-
sure-restr-to-subalg subalg)
   hence emeasure (restr-to-subalg M F) X > 0 using \langle \neg (emeasure M X) = 0 \rangle
gr-zeroI by auto
    then obtain A where A: A \in sets (restr-to-subalg M F) A \subseteq X emeasure
(restr-to-subalg M F) A > 0 emeasure (restr-to-subalg M F) A < \infty
   using sigma-fin-subalg by (metis emeasure-notin-sets ennreal-0 infinity-ennreal-def
le-less-linear\ neq-top-trans\ not-gr-zero\ order-refl\ sigma-finite-measure.approx-PInf-emeasure-with-finite)
   hence [simp]: A \in sets \ F  using subalq \ sets-restr-to-subalq by blast
   hence A-in-sets-M[simp]: A \in sets \ M using sets-restr-to-subalq subalq subal-
qebra-def by blast
    have [simp]: set-integrable M A (\lambda x. c) using A subalg by (auto simp add:
set-integrable-def emeasure-restr-to-subalg)
   have [simp]: set-integrable M A f unfolding set-integrable-def by (rule inte-
grable-mult-indicator, auto\ simp\ add:\ assms(1))
   have AE x in M. indicator A x *_R c = indicator A x *_R f x
   proof (rule integral-eq-mono-AE-eq-AE)
   have (\int x \in A. \ c \ \partial M) \le (\int x \in A. \ fx \ \partial M) using assms(2) by (intro\ set\text{-integral-mono-}AE\text{-}banach)
auto
     moreover
     {
          have (\int x \in A. \ f \ x \ \partial M) = (\int x \in A. \ cond\text{-}exp \ M \ F \ f \ x \ \partial M) by (rule
cond-exp-set-integral, auto simp add: assms)
    also have ... \leq (\int x \in A. \ c \ \partial M) using A by (auto intro!: set-integral-mono-banach
simp \ add: X-def)
       finally have (\int x \in A. \int x \, \partial M) \leq (\int x \in A. \, c \, \partial M) by simp
     ultimately show LINT x|M. indicator A \times_R c = LINT \times_R M. indicator A
x *_R f x unfolding set-lebesgue-integral-def by simp
    show AE x in M. indicator A x *_{B} c < indicator A x *_{B} f x using assms by
(auto simp add: X-def indicator-def)
   qed (auto simp add: set-integrable-def[symmetric])
   hence AE \ x \in A \ in \ M. \ c = f \ x \ by \ auto
   hence AE \ x \in A \ in \ M. False using assms(2) by auto
   hence A \in null-sets M using AE-iff-null-sets A-in-sets-M by metis
    thus False using A(3) by (simp add: emeasure-restr-to-subalg null-setsD1
subalg)
 thus ?thesis using AE-iff-null-sets[OF X-in-M] unfolding X-def by auto
qed
corollary cond-exp-less-c:
  fixes f: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
```

```
dered-real-vector}
 assumes integrable M f AE x in M. f x < c
 shows AE x in M. cond-exp M F f x < c
proof -
  have AE x in M. cond-exp M F f x = - cond-exp M F (\lambda x. - f x) x using
cond-exp-uminus[OF assms(1)] by auto
 moreover have AE x in M. cond-exp M F (\lambda x. - f x) x > -c using assms
by (intro cond-exp-gr-c) auto
 ultimately show ?thesis by (force simp add: minus-less-iff)
qed
lemma cond-exp-mono-strict:
  fixes f::'a \Rightarrow 'b:: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes integrable M f integrable M g AE x in M. f x < g x
 shows AE x in M. cond-exp M F f x < cond-exp M F q x
 using cond-exp-less-c[OF Bochner-Integration.integrable-diff, OF assms(1,2), of
\theta
      cond-exp-diff[OF assms(1,2)] assms(3) by auto
lemma cond-exp-ge-c:
  fixes f: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes [measurable]: integrable M f
     and AE x in M. f x \geq c
 shows AE x in M. cond-exp M F f x \ge c
proof -
 let ?F = restr-to-subalq M F
 interpret sigma-finite-measure restr-to-subalg M F using sigma-fin-subalg by
auto
   fix A assume asm: A \in sets ?F 0 < measure ?F A
  have [simp]: sets ?F = sets\ F\ measure\ ?F\ A = measure\ M\ A\ using\ asm\ by\ (auto
simp\ add: measure-def\ sets-restr-to-subalg[OF\ subalg] emeasure-restr-to-subalg[OF\ subalg]
subalg])
   have M-A: emeasure M A < \infty using measure-zero-top asm by (force simp
add: top.not-eq-extremum)
   hence F-A: emeasure ?F A < \infty using asm(1) emeasure-restr-to-subalg subalg
by fastforce
    have set-lebesgue-integral M A (\lambda-. c) \leq set-lebesgue-integral M A f using
assms asm M-A subalg by (intro set-integral-mono-AE-banach, auto simp add:
set-integrable-def integrable-mult-indicator subalgebra-def sets-restr-to-subalg)
  also have ... = set-lebesque-integral M A (cond-exp M F f) using cond-exp-set-integral [OF
assms(1)] asm by auto
  also have \dots = set-lebesgue-integral ?F A (cond-exp M F f) unfolding set-lebesgue-integral-def
using asm borel-measurable-cond-exp by (intro integral-subalgebra2 OF subalg, sym-
metric, simp)
  finally have (1 / measure ?FA) *_R set-lebesgue-integral ?FA (cond-exp M F f)
\in \{c..\} using asm subalg M-A by (auto simp add: set-integral-const subalgebra-def
```

```
intro!: pos-divideR-le-eq[THEN iffD1])
  thus ?thesis using AE-restr-to-subalg[OF subalg] averaging-theorem[OF inte-
qrable-in-subalg closed-atLeast, OF subalg borel-measurable-cond-exp integrable-cond-exp
by auto
qed
corollary cond-exp-le-c:
  fixes f: 'a \Rightarrow 'b: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes integrable M f
     and AE x in M. f x \leq c
 shows AE x in M. cond-exp M F f x \leq c
proof -
  have AE x in M. cond-exp M F f x = - cond-exp M F (\lambda x. - f x) x using
cond-exp-uminus[OF assms(1)] by force
  moreover have AE x in M. cond-exp M F (\lambda x. - f x) x \ge -c using assms
by (intro cond-exp-ge-c) auto
 ultimately show ?thesis by (force simp add: minus-le-iff)
qed
corollary cond-exp-mono:
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes integrable M f integrable M g AE x in M. f x \leq g x
 shows AE \ x \ in \ M. \ cond\text{-}exp \ M \ F \ f \ x \leq cond\text{-}exp \ M \ F \ g \ x
  using cond-exp-le-c[OF Bochner-Integration.integrable-diff, OF assms(1,2), of
\theta
       cond-exp-diff[OF assms(1,2)] assms(3) by auto
corollary cond-exp-min:
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes integrable M f integrable M g
 shows AE \xi in M. cond-exp M F (\lambda x. min (f x) (g x)) \xi \leq min (cond-exp <math>M F
f \ \xi) (cond-exp M F q \ \xi)
proof -
 have AE \ \xi \ in \ M. \ cond-exp \ M \ F \ (\lambda x. \ min \ (f \ x) \ (g \ x)) \ \xi \leq cond-exp \ M \ F \ f \ \xi \ by
(intro cond-exp-mono integrable-min assms, simp)
  moreover have AE \xi in M. cond\text{-}exp M F (\lambda x. min (f x) (g x)) \xi \leq cond\text{-}exp
M F g \xi by (intro cond-exp-mono integrable-min assms, simp)
  ultimately show AE \xi in M. cond-exp M F (\lambda x. min (f x) (g x)) \xi \leq min
(cond\text{-}exp\ M\ F\ f\ \xi)\ (cond\text{-}exp\ M\ F\ g\ \xi)\ \mathbf{by}\ fastforce
qed
corollary cond-exp-max:
  fixes f: 'a \Rightarrow 'b: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes integrable M f integrable M g
```

```
shows AE \ \xi in M. cond-exp M \ F \ (\lambda x. \ max \ (f \ x) \ (g \ x)) \ \xi \ge max \ (cond-exp \ M \ F
f \xi) (cond-exp M F g \xi)
proof -
 have AE \ \xi \ in \ M. \ cond-exp \ M \ F \ (\lambda x. \ max \ (f \ x) \ (g \ x)) \ \xi \geq cond-exp \ M \ F \ f \ \xi \ by
(intro cond-exp-mono integrable-max assms, simp)
 moreover have AE \xi in M. cond-exp M F (\lambda x. max (f x) (g x)) \xi \geq cond-exp
M F g \xi by (intro cond-exp-mono integrable-max assms, simp)
  ultimately show AE \ \xi \ in \ M. \ cond-exp \ M \ F \ (\lambda x. \ max \ (f \ x) \ (g \ x)) \ \xi \geq max
(cond\text{-}exp\ M\ F\ f\ \xi)\ (cond\text{-}exp\ M\ F\ g\ \xi)\ \mathbf{by}\ fastforce
qed
corollary cond-exp-inf:
  fixes f::'a \Rightarrow 'b:: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector, lattice}
 assumes integrable M f integrable M g
 shows AE \notin in M. cond-exp M F (\lambda x. inf (f x) (g x)) \notin inf (cond-exp M F f
\xi) (cond-exp M F g \xi)
 unfolding inf-min using assms by (rule cond-exp-min)
corollary cond-exp-sup:
  fixes f: 'a \Rightarrow 'b: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector, lattice}
 assumes integrable M f integrable M g
 shows AE \xi in M. cond-exp M F (\lambda x. sup (f x) (g x)) \xi \ge sup (cond-exp <math>M F f
\xi) (cond-exp M F q \xi)
  unfolding sup-max using assms by (rule cond-exp-max)
end
4.4
        Probability Spaces
lemma (in prob-space) sigma-finite-subalgebra-restr-to-subalg:
 assumes subalgebra M F
 shows sigma-finite-subalgebra M F
proof (intro sigma-finite-subalgebra.intro)
 interpret\ F:\ prob\ -space\ restr-to\ -subalg\ M\ F\ using\ assms\ prob\ -space\ -restr-to\ -subalg
prob-space-axioms by blast
 show sigma-finite-measure (restr-to-subalq MF) by (rule F.sigma-finite-measure-axioms)
qed (rule assms)
lemma (in prob-space) cond-exp-trivial:
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
 assumes integrable M f
 shows AE x in M. cond-exp M (sigma (space M) \{\}) f x = expectation f
 interpret sigma-finite-subalgebra M sigma (space M) {} by (auto intro: sigma-finite-subalgebra-restr-to-subalgebra
simp add: subalgebra-def sigma-sets-empty-eq)
 show ?thesis using assms by (intro cond-exp-charact) (auto simp add: sigma-sets-empty-eq
set-lebesque-integral-def prob-space conq: Bochner-Integration.integral-conq)
```

qed

The following lemma shows that independent sigma algebras don't matter for the conditional expectation.

```
lemma (in prob-space) cond-exp-indep-subalgebra:
 fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, real\text{-}normed\text{-}field}\}
 assumes subalgebra: subalgebra M F subalgebra M G
     and independent: indep-set G (sigma (space M) (F \cup vimage-algebra (space
M) f borel)
 assumes [measurable]: integrable M f
 shows AE \times in M. cond-exp M (sigma (space M) (F \cup G)) f \times cond-exp M F
f x
proof -
 interpret Un-sigma: sigma-finite-subalgebra M sigma (space M) (F \cup G) using
assms(1,2) by (auto intro!: sigma-finite-subalgebra-restr-to-subalg sets.sigma-sets-subset
simp add: subalgebra-def space-measure-of-conv sets-measure-of-conv)
 interpret sigma-finite-subalgebra MF using assms by (auto intro: sigma-finite-subalgebra-restr-to-subalg)
 {
   \mathbf{fix} \ A
   assume asm: A \in sigma \ (space \ M) \ \{a \cap b \mid a \ b. \ a \in F \land b \in G\}
    have in-events: sigma-sets (space M) \{a \cap b \mid a \ b. \ a \in sets \ F \land b \in sets \}
G \subseteq events  using subalgebra  by (intro sets.sigma-sets-subset, auto simp add:
subalgebra-def)
   have Int-stable \{a \cap b \mid a \ b. \ a \in F \land b \in G\}
   proof -
       fix af bf ag bg
       assume F: af \in F \ bf \in F \ and \ G: ag \in G \ bg \in G
       have af \cap bf \in F by (intro sets.Int F)
       moreover have ag \cap bg \in G by (intro sets.Int G)
       ultimately have \exists a \ b. \ af \cap ag \cap (bf \cap bg) = a \cap b \wedge a \in sets \ F \wedge b \in
sets G by (metis inf-assoc inf-left-commute)
     }
     thus ?thesis by (force intro!: Int-stableI)
    moreover have \{a \cap b \mid a \ b. \ a \in F \land b \in G\} \subseteq Pow \ (space \ M) using
subalgebra by (force simp add: subalgebra-def dest: sets.sets-into-space)
   moreover have A \in sigma\text{-}sets (space M) \{a \cap b \mid a \ b. \ a \in F \land b \in G\} using
calculation asm by force
     ultimately have set-lebesgue-integral M A f = set-lebesgue-integral M A
(cond\text{-}exp\ M\ F\ f)
   proof (induction rule: sigma-sets-induct-disjoint)
     case (basic A)
     then obtain a b where A: A = a \cap b a \in F b \in G by blast
     hence events[measurable]: a \in events b \in events using subalgebra by (auto
simp add: subalgebra-def)
```

have [simp]: sigma-sets (space M) {indicator $b - A \cap space M \mid A. A \in borel}$

 $\subset G$

 $\begin{tabular}{ll} \textbf{using} \ borel-measurable-indicator} [OF\ A(3),\ THEN\ measurable-sets] \ sets.top \\ subalgebra \end{tabular}$

by (intro sets.sigma-sets-subset') (fastforce simp add: subalgebra-def)+

have Un-in-sigma: $F \cup vimage$ -algebra (space M) f borel $\subseteq sigma$ (space M) ($F \cup vimage$ -algebra (space M) f borel) by (metis equality E le-sup I sets.space-closed sigma-le-sets space-vimage-algebra subalg subalgebra-def)

have [intro]: indep-var borel (indicator b) borel ($\lambda \omega$. indicator a $\omega *_R f \omega$) proof -

have [simp]: sigma-sets (space M) {($\lambda\omega$. indicator a $\omega *_R f \omega$) – ' $A \cap$ space M |A. $A \in borel$ } $\subseteq sigma$ (space M) ($F \cup vimage-algebra$ (space M) f borel) proof –

have *: $(\lambda \omega$. indicator a $\omega *_R f \omega) \in borel$ -measurable (sigma (space M) $(F \cup vimage$ -algebra (space M) f borel))

using borel-measurable-indicator [OF A(2), THEN measurable-sets, OF borel-open] subalgebra

 $\mathbf{by}\ (intro\ borel-measurable\text{-}scaleR\ borel-measurableI\ Un-in\text{-}sigma[THEN\ subsetD])}$

(auto simp add: space-measure-of-conv subalgebra-def sets-vimage-algebra2) thus ?thesis using measurable-sets[OF *] by (intro sets.sigma-sets-subset', auto simp add: space-measure-of-conv)

qed

have indep-set (sigma-sets (space M) {indicator b - 'A \cap space M | A. $A \in borel$ }) (sigma-sets (space M) {($\lambda \omega$. indicator $a \omega *_R f \omega$) - ' $A \cap space M \mid A$. $A \in borel$ })

using independent unfolding indep-set-def by (rule indep-sets-mono-sets, auto split: bool.split)

thus ?thesis by (subst indep-var-eq, auto intro!: borel-measurable-scaleR) qed

have [intro]: indep-var borel (indicator b) borel ($\lambda\omega$. indicat-real a $\omega*_R$ cond-exp M F f ω)

proof -

have [simp]: sigma-sets $(space\ M)\ \{(\lambda\omega.\ indicator\ a\ \omega*_R\ cond-exp\ M\ F\ f\ \omega)$ - ' $A\cap space\ M\ |A.\ A\in borel\}\subseteq sigma\ (space\ M)\ (F\cup vimage-algebra\ (space\ M)\ f\ borel)$

proof -

have *: $(\lambda \omega. indicator \ a \ \omega *_R cond-exp \ M \ F \ f \ \omega) \in borel-measurable (sigma (space M) (F \cup vimage-algebra (space M) f borel))$

 $\begin{tabular}{ll} \textbf{using} \begin{tabular}{ll} borel-measurable-indicator[OF\ A(2),\ THEN\ measurable-sets,\ OF\ borel-open]\ subalgebra \end{tabular}$

 $borel-measurable\text{-}cond\text{-}exp[\mathit{THEN}\ measurable\text{-}sets,\ OF\ borel\text{-}open,\ of\ -\ M\ F\ f]$

 $\mathbf{by}\ (intro\ borel-measurable\text{-}scaleR\ borel-measurableI\ Un\text{-}in\text{-}sigma[THEN\ subsetD]})$

(auto simp add: space-measure-of-conv subalgebra-def)

thus ?thesis using measurable-sets[OF *] by (intro sets.sigma-sets-subset',

```
qed
              have indep-set (sigma-sets (space M) {indicator b - A \cap space M \mid A. A \in A \cap space M \mid A. A \in A \cap space M \mid A. A \in A \cap space M \mid A \cap space
borel}) (sigma-sets (space M) \{(\lambda \omega. indicator \ a \ \omega *_R \ cond-exp \ M \ F \ f \ \omega) - `A \cap A \}
space\ M\ |A.\ A \in borel\}
                 using independent unfolding indep-set-def by (rule indep-sets-mono-sets,
auto split: bool.split)
                thus ?thesis by (subst indep-var-eq, auto intro!: borel-measurable-scaleR)
            qed
            have set-lebesgue-integral M A f = (LINT \ x|M. \ indicator \ b \ x * (indicator \ a
x *_R f x)
               unfolding set-lebesque-integral-def A indicator-inter-arith
           by (intro Bochner-Integration.integral-cong, auto simp add: scaleR-scaleR[symmetric]
indicator-times-eq-if(1)
            also have ... = (LINT \ x|M. \ indicator \ b \ x) * (LINT \ x|M. \ indicator \ a \ x*_B \ f
x)
               by (intro indep-var-lebesgue-integral
                                     Bochner-Integration.integrable-bound[OF integrable-const[of 1 :: 'b]]
borel-measurable-indicator]
                                          integrable-mult-indicator[OF - assms(4)], blast) (auto simp \ add:
indicator-def)
                also have ... = (LINT \ x|M. \ indicator \ b \ x) * (LINT \ x|M. \ indicator \ a \ x*_R
cond-exp M F f x)
              using cond-exp-set-integral [OF assms(4) A(2)] unfolding set-lebesque-integral-def
                also have ... = (LINT x|M. indicator b x * (indicator \ a \ x *_R \ cond-exp \ M
Ff(x)
               by (intro indep-var-lebesgue-integral[symmetric]
                                     Bochner-Integration.integrable-bound[OF integrable-const[of 1 :: 'b]
borel-measurable-indicator]
                                integrable-mult-indicator[OF - integrable-cond-exp], blast) (auto simp
add: indicator-def)
            also have ... = set-lebesgue-integral M A (cond-exp M F f)
               unfolding set-lebesgue-integral-def A indicator-inter-arith
           by (intro Bochner-Integration.integral-cong, auto simp add: scaleR-scaleR[symmetric]
indicator-times-eq-if(1)
            finally show ?case.
       next
            case empty
            then show ?case unfolding set-lebesgue-integral-def by simp
       next
            case (compl A)
        have A-in-space: A \subseteq space \ M using compl using in-events sets.sets-into-space
         have set-lebesgue-integral M (space M-A) f=set-lebesgue-integral M (space
M-A\cup A) f- set-lebesgue-integral M A f
               using compl(1) in-events
               by (subst set-integral-Un[of space M - A A], blast)
```

auto simp add: space-measure-of-conv)

```
(simp | intro integrable-mult-indicator folded set-integrable-def, OF -
assms(4)], fast)+
    also have ... = set-lebesgue-integral M (space M - A \cup A) (cond-exp M F f)
- set-lebesque-integral M A (cond-exp M F f)
     using cond-exp-set-integral [OF assms(4) sets.top] compl subalgebra by (simp
add: subalgebra-def Un-absorb2[OF A-in-space])
     also have \dots = set-lebesgue-integral M (space M-A) (cond-exp M F f)
       using compl(1) in-events
      by (subst set-integral-Un[of space M - A A], blast)
             (simp \mid intro integrable-mult-indicator [folded set-integrable-def, OF -
integrable-cond-exp[, fast)+
     finally show ?case.
   \mathbf{next}
     case (union A)
    have set-lebesgue-integral M (() (range A)) f = (\sum i. set-lebesgue-integral M)
(A \ i) \ f)
       using union in-events
     by (intro lebesgue-integral-countable-add) (auto simp add: disjoint-family-onD
intro!: integrable-mult-indicator[folded\ set-integrable-def,\ OF\ -\ assms(4)])
     also have ... = (\sum i. set-lebesgue-integral M (A i) (cond-exp M F f)) using
union by presburger
     also have ... = set-lebesgue-integral M (\bigcup (range \ A)) (cond-exp M \ F \ f)
       using union in-events
       by (intro lebesque-integral-countable-add[symmetric]) (auto simp add: dis-
joint-family-onD introl: integrable-mult-indicator[folded set-integrable-def, OF - in-
tegrable-cond-exp])
     finally show ?case.
   qed
 }
 moreover have sigma\ (space\ M)\ \{a\cap b\mid a\ b.\ a\in F\land b\in G\}=sigma\ (space\ M)
M) (F \cup G)
 proof -
   have sigma-sets (space M) \{a \cap b \mid a \ b. \ a \in sets \ F \land b \in sets \ G\} = sigma-sets
(space\ M)\ (sets\ F\cup sets\ G)
   proof -
      fix a b assume asm: a \in F b \in G
      hence a \cap b \in sigma\text{-}sets (space M) (F \cup G) using subalgebra unfolding
Int-range-binary by (intro sigma-sets-Inter[OF - binary-in-sigma-sets]) (force simp
add: subalgebra-def dest: sets.sets-into-space)+
     }
     moreover
     {
      \mathbf{fix} \ a
      assume a \in sets F
      hence a \in sigma\text{-}sets (space M) \{a \cap b \mid a \text{ b. } a \in sets F \land b \in sets G\}
        using subalgebra sets.top[of G] sets.sets-into-space[of - F]
        by (intro sigma-sets.Basic, auto simp add: subalgebra-def)
     }
```

```
moreover
      fix a assume a \in sets \ F \lor a \in sets \ G \ a \notin sets \ F
      hence a \in sets G by blast
      hence a \in sigma\text{-}sets (space M) \{a \cap b \mid a b.\ a \in sets\ F \land b \in sets\ G\}
        using subalgebra sets.top[of F] sets.sets-into-space[of - G]
        by (intro sigma-sets. Basic, auto simp add: subalgebra-def)
     ultimately show ?thesis by (intro sigma-sets-eqI) auto
   qed
   thus ?thesis using subalgebra by (intro sigma-eqI) (force simp add: subalge-
bra-def dest: sets.sets-into-space)+
 moreover have (cond\text{-}exp\ M\ F\ f) \in borel\text{-}measurable\ (sigma\ (space\ M)\ (sets\ F\ )
\cup sets G))
  proof -
    have F \subseteq sigma \ (space \ M) \ (F \cup G) by (metis Un-least Un-upper1 mea-
sure-of-of-measure sets.space-closed sets-measure-of sigma-sets-subseteq subalg sub-
algebra(2) subalgebra-def)
  thus ?thesis using borel-measurable-cond-exp[THEN measurable-sets, OF borel-open,
of - M F f subalgebra by (intro borel-measurable I, force simp only: space-measure-of-conv
subalgebra-def)
  qed
 ultimately show ?thesis using assms(4) integrable-cond-exp by (intro Un-sigma.cond-exp-charact)
presburger +
qed
lemma (in prob-space) cond-exp-indep:
 fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, real\text{-}normed\text{-}field}\}
 assumes subalgebra: subalgebra M F
     and independent: indep-set \ F \ (vimage-algebra \ (space \ M) \ f \ borel)
     and integrable: integrable M f
 shows AE x in M. cond-exp M F f x = expectation f
proof -
 have indep\text{-}set\ F\ (sigma\ (space\ M)\ (sigma\ (space\ M)\ \{\}\cup(vimage\text{-}algebra\ (space\ M)\ \{\})
M) f borel)))
   using independent unfolding indep-set-def
   by (rule indep-sets-mono-sets, simp add: bool.split)
    (metis bot.extremum dual-order.refl sets.sets-measure-of-eq sets.sigma-sets-subset'
sets-vimage-algebra-space space-vimage-algebra sup.absorb-iff2)
 hence cond-exp-indep: AE x in M. cond-exp M (sigma (space M) (sigma (space
M) {} \cup F)) f x = expectation <math>f
    using cond-exp-indep-subalgebra [OF - subalgebra - integrable, of sigma (space
M) {}| cond-exp-trivial[OF integrable]
   by (auto simp add: subalgebra-def sigma-sets-empty-eq)
  have sets (sigma (space M) (sigma (space M) \{\} \cup F\}) = F
   using subalgebra\ sets.top[of\ F] unfolding subalgebra-def
   by (simp add: sigma-sets-empty-eq, subst insert-absorb[of space M F], blast)
      (metis insert-absorb[OF sets.empty-sets] sets.sets-measure-of-eq)
```

```
hence AE\ x in M. cond-exp\ M (sigma\ (space\ M) (sigma\ (space\ M) {} \cup\ F)) f x=cond-exp\ M\ F\ f\ x by (rule\ cond-exp-sets-cong) thus ?thesis\ using\ cond-exp-indep\ by\ force qed end theory Filtered-Measure\ imports\ HOL-Probability. Conditional-Expectation begin
```

5 Filtered Measure Spaces

5.1 Filtered Measure

```
locale filtered-measure =
 fixes M F and t_0 :: \{second\text{-}countable\text{-}topology, order\text{-}topology, t2\text{-}space}\}
 assumes subalgebras: \bigwedge i. t_0 \leq i \Longrightarrow subalgebra\ M\ (F\ i)
     and sets-F-mono: \bigwedge i \ j. t_0 \le i \Longrightarrow i \le j \Longrightarrow sets \ (F \ i) \le sets \ (F \ j)
begin
lemma space-F[simp]:
 assumes t_0 \leq i
 shows space (F i) = space M
 using subalgebras assms by (simp add: subalgebra-def)
lemma subalgebra-F[intro]:
 assumes t_0 \leq i \ i \leq j
 shows subalgebra (F j) (F i)
 unfolding subalgebra-def using assms by (simp add: sets-F-mono)
{f lemma}\ borel-measurable-mono:
  assumes t_0 \leq i \ i \leq j
 shows borel-measurable (F i) \subseteq borel-measurable (F j)
 unfolding subset-iff by (metis assms subalgebra-F measurable-from-subalg)
end
locale linearly-filtered-measure = filtered-measure M F t_0 for M and F :: - ::
\{linorder\text{-}topology\} \Rightarrow - \text{ and } t_0
locale nat-filtered-measure = linearly-filtered-measure M F \theta for M and F :: nat
locale real-filtered-measure = linearly-filtered-measure M F \theta for M and F :: real
```

5.2 Sigma Finite Filtered Measure

The locale presented here is a generalization of the *sigma-finite-subalgebra* for a particular filtration.

```
locale \ sigma-finite-filtered-measure = filtered-measure +
 assumes sigma-finite-initial: sigma-finite-subalgebra M (F t_0)
\mathbf{lemma} \ (\mathbf{in} \ sigma-finite-filtered-measure}) \ sigma-finite-subalgebra-F[intro]:
 assumes t_0 \leq i
 shows sigma-finite-subalgebra M (F i)
 \textbf{using} \ assms \ \textbf{by} \ (metis \ dual-order. refl \ sets-F-mono \ sigma-finite-initial \ sigma-finite-subalgebra. nested-subalg-is
subalgebras subalgebra-def)
locale nat-sigma-finite-filtered-measure = sigma-finite-filtered-measure M F 0 ::
nat for M F
locale real-sigma-finite-filtered-measure = sigma-finite-filtered-measure M F 0 ::
real for M F
sublocale nat-sigma-finite-filtered-measure \subseteq sigma-finite-subalgebra M F i by
sublocale real-sigma-finite-filtered-measure \subseteq sigma-finite-subalgebra M F |i| by
fastforce
5.3
       Finite Filtered Measure
locale\ finite-filtered-measure\ =\ filtered-measure\ +\ finite-measure
sublocale finite-filtered-measure \subseteq sigma-finite-filtered-measure
 using subalgebras by (unfold-locales, blast, meson dual-order.reft finite-measure-axioms
finite-measure-def finite-measure-restr-to-subalg sigma-finite-measure.sigma-finite-countable)
locale nat-finite-filtered-measure = finite-filtered-measure M F \theta :: nat for M F
locale real-finite-filtered-measure = finite-filtered-measure M F \theta :: real for M F
\textbf{sublocale} \ \textit{nat-finite-filtered-measure} \subseteq \textit{nat-sigma-finite-filtered-measure} \ ..
\mathbf{sublocale} real-finite-filtered-measure \subseteq real-sigma-finite-filtered-measure ...
```

5.4 Constant Filtration

```
lemma filtered-measure-constant-filtration:

assumes subalgebra M F

shows filtered-measure M (\lambda-. F) t_0

using assms by (unfold-locales) blast+

sublocale sigma-finite-subalgebra \subseteq constant-filtration: sigma-finite-filtered-measure

M \lambda- :: 't :: {second-countable-topology, linorder-topology}. F t_0

using subalg by (unfold-locales) blast+

lemma (in finite-measure) filtered-measure-constant-filtration:
```

```
assumes subalgebra M F shows finite-filtered-measure M (\lambda-. F) t_0 using assms by (unfold-locales) blast+
```

end

 ${\bf theory}\ Stochastic-Process\\ {\bf imports}\ Filtered-Measure\ Measure-Space-Supplement\ HOL-Probability. Independent-Family\ {\bf begin}$

6 Stochastic Processes

6.1 Stochastic Process

A stochastic process is a collection of random variables, indexed by a type 'b.

```
locale stochastic-process =
 fixes M t_0 and X :: 'b :: \{second\text{-}countable\text{-}topology, order\text{-}topology, t2\text{-}space}\} \Rightarrow
'a \Rightarrow 'c :: \{second\text{-}countable\text{-}topology, banach}\}
 assumes random-variable[measurable]: \bigwedge i. t_0 \leq i \Longrightarrow X i \in borel-measurable M
begin
definition left-continuous where left-continuous = (AE \ \xi \ in \ M. \ \forall \ t. \ continuous
(at\text{-left }t)\ (\lambda i.\ X\ i\ \xi))
definition right-continuous where right-continuous = (AE \xi \text{ in } M. \forall t. \text{ continuous})
(at\text{-}right\ t)\ (\lambda i.\ X\ i\ \xi))
end
\mathbf{locale}\ nat\text{-}stochastic\text{-}process = stochastic\text{-}process\ M\ 0 :: nat\ X\ \mathbf{for}\ M\ X
locale \ real-stochastic-process = stochastic-process M 0 :: real X for M X
lemma stochastic-process-const-fun:
  assumes f \in borel-measurable M
 shows stochastic-process M t_0 (\lambda-. f) using assms by (unfold-locales)
lemma stochastic-process-const:
  shows stochastic-process M t_0 (\lambda i -. c i) by (unfold-locales) simp
context stochastic-process
begin
lemma compose-stochastic:
 assumes \bigwedge i. t_0 \leq i \Longrightarrow f i \in borel-measurable borel
  shows stochastic-process M t_0 (\lambda i \ \xi. (f \ i) (X \ i \ \xi))
  by (unfold-locales) (intro measurable-compose[OF random-variable assms])
```

```
lemma norm-stochastic: stochastic-process M t_0 (\lambda i \ \xi. norm (X \ i \ \xi)) by (fastforce
intro: compose-stochastic)
lemma scaleR-right-stochastic:
  assumes stochastic-process\ M\ t_0\ Y
 shows stochastic-process M t_0 (\lambda i \ \xi. (Y \ i \ \xi) *_R (X \ i \ \xi))
 \mathbf{using}\ stochastic-process.random-variable[OF\ assms]\ random-variable\ \mathbf{by}\ (unfold-locales)
\mathbf{lemma}\ scaleR-right-const-fun-stochastic:
  assumes f \in borel-measurable M
  shows stochastic-process M t_0 (\lambda i \ \xi. \ f \ \xi *_R (X \ i \ \xi))
  \mathbf{by}\ (unfold\text{-}locales)\ (intro\ borel\text{-}measurable\text{-}scaleR\ assms\ random\text{-}variable)
lemma scaleR-right-const-stochastic: stochastic-process M t_0 (\lambda i \ \xi. \ c \ i *_R (X \ i \ \xi))
  by (unfold-locales) simp
lemma add-stochastic:
  assumes stochastic-process\ M\ t_0\ Y
 shows stochastic-process M t_0 (\lambda i \xi. X i \xi + Y i \xi)
 \mathbf{using}\ stochastic\text{-}process.random\text{-}variable[OF\ assms]\ random\text{-}variable\ \mathbf{by}\ (unfold\text{-}locales)
simp
lemma diff-stochastic:
  assumes stochastic-process\ M\ t_0\ Y
  shows stochastic-process M t_0 (\lambda i \ \xi. X \ i \ \xi - Y \ i \ \xi)
 using stochastic-process.random-variable [OF assms] random-variable by (unfold-locales)
simp
lemma uminus-stochastic: stochastic-process M t_0 (-X) using scaleR-right-const-stochastic [of
\lambda-. -1] by (simp add: fun-Compl-def)
lemma partial-sum-stochastic: stochastic-process M t_0 (\lambda n \xi. \sum i \in \{t_0..n\}. X i \xi)
by (unfold-locales) simp
lemma partial-sum'-stochastic: stochastic-process M t_0 (\lambda n \xi. \sum i \in \{t_0... < n\}). X i
\xi) by (unfold-locales) simp
end
lemma stochastic-process-sum:
  assumes \bigwedge i. i \in I \Longrightarrow stochastic-process M t_0 (X i)
  shows stochastic-process M t_0 (\lambda k \xi. \sum i \in I. X i k \xi) using assms[THEN
stochastic-process.random-variable] by (unfold-locales, auto)
```

6.1.1 Natural Filtration

The natural filtration induced by a stochastic process X is the filtration generated by all events involving the process up to the time index t, i.e. Σ

```
t = \sigma \{X \mid s \mid s \leq t\}.
definition natural-filtration :: 'a measure \Rightarrow 'b \Rightarrow ('b \Rightarrow 'a \Rightarrow 'c :: topologi-
cal\text{-}space) \Rightarrow b' : \{second\text{-}countable\text{-}topology, order\text{-}topology}\} \Rightarrow a' measure where
  natural-filtration M t_0 Y = (\lambda t. family-vimage-algebra (space M) { Y i | i. i \in
\{t_0..t\}\} borel)
abbreviation nat-natural-filtration \equiv \lambda M. natural-filtration M (0 :: nat)
abbreviation real-natural-filtration \equiv \lambda M. natural-filtration M (0 :: real)
lemma space-natural-filtration[simp]: space (natural-filtration M t_0 X t) = space
M unfolding natural-filtration-def space-family-vimage-algebra ...
lemma sets-natural-filtration: sets (natural-filtration M t_0 X t) = sigma-sets (space
M) (\bigcup i \in \{t_0...t\}. \{X \ i - `A \cap space M \mid A. A \in borel\})
 unfolding natural-filtration-def sets-family-vimage-algebra by (intro sigma-sets-eqI)
blast+
lemma sets-natural-filtration':
 assumes borel = sigma\ UNIV\ S
 shows sets (natural-filtration M t_0 X t) = sigma-sets (space M) (\bigcup i \in \{t_0..t\}. \{X\}
i - A \cap space M \mid A. A \in S\}
proof (subst sets-natural-filtration, intro sigma-sets-eqI, clarify)
  fix i and A :: 'a set assume asm: i \in \{t_0..t\} A \in sets borel
  hence A \in sigma\text{-}sets \ UNIV \ S \ \mathbf{unfolding} \ assms \ \mathbf{by} \ simp
  thus X i - A \cap space M \in sigma-sets (space M) (\bigcup i \in \{t_0..t\}. \{X i - A \cap A \cap A \cap A\})
space M \mid A. A \in S \}
 proof (induction)
   case (Compl\ a)
   have X i - (UNIV - a) \cap space M = space M - (X i - a \cap space M) by
blast
   then show ?case using Compl(2)[THEN sigma-sets.Compl] by presburger
 next
   case (Union \ a)
   have X \ i - `(\bigcup (range \ a) \cap space \ M = \bigcup (range \ (\lambda j. \ X \ i - `a \ j \cap space \ M))
   then show ?case using Union(2)[THEN sigma-sets.Union] by presburger
 qed (auto intro: asm sigma-sets.Empty)
qed (intro sigma-sets.Basic, force simp add: assms)
lemma sets-natural-filtration-open:
  sets (natural-filtration M t_0 X t) = sigma-sets (space M) (\bigcup i \in \{t_0..t\}. \{X i - i\}
A \cap space M \mid A. open A\}
 using sets-natural-filtration' by (force simp only: borel-def mem-Collect-eq)
lemma sets-natural-filtration-oi:
 sets (natural-filtration M t_0 X t) = sigma-sets (space M) (\bigcup i \in \{t_0..t\}. {X i - 'A
```

 \cap space $M \mid A :: - :: \{linorder-topology, second-countable-topology\} set. <math>A \in range$

 $greaterThan\})$

by (rule sets-natural-filtration'[OF borel-Ioi])

```
lessThan\})
   by (rule sets-natural-filtration'[OF borel-Iio])
lemma sets-natural-filtration-ci:
    sets (natural-filtration M t_0 X t) = sigma-sets (space M) (\bigcup i \in \{t_0..t\}. {X i - '
A \cap space M \mid A :: real set. A \in range atLeast\}
   by (rule sets-natural-filtration'[OF borel-Ici])
context stochastic-process
begin
lemma subalgebra-natural-filtration:
   shows subalgebra M (natural-filtration M t_0 X i)
    unfolding subalgebra-def using measurable-family-iff-sets by (force simp add:
natural-filtration-def)
lemma filtered-measure-natural-filtration:
   shows filtered-measure M (natural-filtration M t_0 X) t_0
    by (unfold-locales) (intro subalgebra-natural-filtration, simp only: sets-natural-filtration,
intro sigma-sets-subseteq, force)
In order to show that the natural filtration constitutes a filtered sigma finite
measure, we need to provide a countable exhausting set in the preimage of
X t_0.
lemma sigma-finite-filtered-measure-natural-filtration:
   assumes exhausting-set: countable A (\bigcup A) = space M \land a. \ a \in A \Longrightarrow emeasure
M \ a \neq \infty \land a. \ a \in A \Longrightarrow \exists \ b \in borel. \ a = X \ t_0 - b \cap space M
       shows sigma-finite-filtered-measure M (natural-filtration M t_0 X) t_0
proof (unfold-locales)
   have A \subseteq sets (restr-to-subala M (natural-filtration M t_0 \times t_0)) using exhaust-
ing-set by (simp add: sets-restr-to-subalg[OF subalgebra-natural-filtration] sets-natural-filtration)
fast
   moreover have [] A = space (restr-to-subalg M (natural-filtration M <math>t_0 X t_0))
unfolding space-restr-to-subalg using exhausting-set by simp
    moreover have \forall a \in A. emeasure (restr-to-subalg M (natural-filtration M t_0 X
t_0) a \neq \infty using calculation(1) exhausting-set(3)
       by (auto simp add: sets-restr-to-subalg[OF subalgebra-natural-filtration] emea-
sure-restr-to-subalq[OF subalgebra-natural-filtration])
  ultimately show \exists A. countable A \land A \subseteq sets (restr-to-subalg M (natural-filtration
M \ t_0 \ X \ t_0) \land \bigcup A = space \ (restr-to-subalg \ M \ (natural-filtration \ M \ t_0 \ X \ t_0)) \land (M \ t_0 \
(\forall a \in A. \ emeasure \ (restr-to-subalg \ M \ (natural-filtration \ M \ t_0 \ X \ t_0)) \ a \neq \infty) using
exhausting-set by blast
  show \bigwedge i j. [t_0 \le i; i \le j] \Longrightarrow sets (natural-filtration M t_0 X i) \subseteq sets (natural-filtration
M t_0 X j) using filtered-measure.subalgebra-F[OF filtered-measure-natural-filtration]
by (simp add: subalgebra-def)
```

sets (natural-filtration M t_0 X t) = sigma-sets (space M) ($\bigcup i \in \{t_0..t\}$. {X i - 'A \cap space $M \mid A :: - :: \{linorder-topology, second-countable-topology\}$ set. $A \in range$

 ${f lemma}\ sets$ -natural-filtration-io:

```
qed (auto intro: subalgebra-natural-filtration)
\mathbf{lemma}\ \mathit{finite-filtered-measure-natural-filtration}:
 assumes finite-measure M
 shows finite-filtered-measure M (natural-filtration M t_0 X) t_0
 using finite-measure.axioms[OF assms] filtered-measure-natural-filtration by in-
tro-locales
end
Filtration generated by independent variables.
lemma (in prob-space) indep-set-natural-filtration:
 assumes t_0 \leq s \leq t \text{ indep-vars } (\lambda \text{-. borel}) X \{t_0..\}
  shows indep-set (natural-filtration M t_0 X s) (vimage-algebra (space M) (X t)
borel)
proof
 \{t_0..s\}\ \{t\})))
   using assms
  by (intro assms(3)[unfolded indep-vars-def, THEN conjunct2, THEN indep-sets-mono])
(auto simp add: case-bool-if)
 thus ?thesis unfolding indep-set-def using assms
   by (intro indep-sets-cong THEN iffD1, OF reft - indep-sets-collect-sigma of \lambda i.
\{X \ i - A \cap space \ M \mid A. \ A \in borel\} \ case-bool \ \{t_0...s\} \ \{t\}\}\}
       (simp add: sets-natural-filtration sets-vimage-algebra split: bool.split, simp,
intro Int-stable I, clarsimp, metis sets. Int vimage-Int Int-commute Int-left-absorb
Int-left-commute, force simp add: disjoint-family-on-def split: bool.split)
qed
6.2
       Adapted Process
We call a collection a stochastic process X adapted if X i is F i-borel-
measurable for all indices i.
locale adapted-process = filtered-measure M F t_0 for M F t_0 and X :: - \Rightarrow - \Rightarrow -
:: \{second\text{-}countable\text{-}topology, banach\} +
  assumes adapted[measurable]: \bigwedge i. t_0 \leq i \Longrightarrow X \ i \in borel-measurable (F i)
begin
lemma adaptedE[elim]:
 assumes \llbracket \bigwedge j \ i. \ t_0 \leq j \Longrightarrow j \leq i \Longrightarrow X \ j \in borel-measurable \ (F \ i) \rrbracket \Longrightarrow P
 using assms using adapted by (metis dual-order trans borel-measurable-subalgebra
sets-F-mono space-F)
lemma adaptedD:
 assumes t_0 \leq j j \leq i
```

shows $X \ j \in borel$ -measurable $(F \ i)$ using assms adapted by meson

end

```
\mathbf{locale} \ \mathit{nat-adapted-process} \ = \ \mathit{adapted-process} \ \mathit{M} \ \mathit{F} \ \mathit{0} \ :: \ \mathit{nat} \ \mathit{X} \ \mathbf{for} \ \mathit{M} \ \mathit{F} \ \mathit{X}
locale real-adapted-process = adapted-process M F 0 :: real X \text{ for } M F X
\mathbf{sublocale}\ \mathit{nat-adapted-process} \subseteq \mathit{nat-filtered-measure}\ ..
sublocale real-adapted-process \subseteq real-filtered-measure ..
lemma (in filtered-measure) adapted-process-const-fun:
  assumes f \in borel-measurable (F t_0)
 shows adapted-process M F t_0 (\lambda-. f)
 using measurable-from-subalg subalgebra-F assms by (unfold-locales) blast
lemma (in filtered-measure) adapted-process-const:
 shows adapted-process M F t_0 (\lambda i - c i) by (unfold-locales) simp
context adapted-process
begin
lemma compose-adapted:
 assumes \bigwedge i. t_0 \leq i \Longrightarrow f i \in borel-measurable borel
 shows adapted-process M F t_0 (\lambda i \xi. (f i) (X i \xi))
 by (unfold-locales) (intro measurable-compose[OF adapted assms])
lemma norm-adapted: adapted-process M F t_0 (\lambda i \xi. norm (X i \xi)) by (fastforce
intro: compose-adapted)
lemma scaleR-right-adapted:
 assumes adapted-process M F t_0 R
 shows adapted-process M F t_0 (\lambda i \xi. (R i \xi) *_R (X i \xi))
 using adapted-process.adapted[OF assms] adapted by (unfold-locales) simp
\mathbf{lemma}\ scaleR-right-const-fun-adapted:
 assumes f \in borel-measurable (F t_0)
 shows adapted-process M F t_0 (\lambda i \xi. f \xi *_R (X i \xi))
 using assms by (fast intro: scaleR-right-adapted adapted-process-const-fun)
lemma scaleR-right-const-adapted: adapted-process M F t_0 (\lambda i \xi. c i *_R (X i \xi))
by (unfold-locales) simp
lemma add-adapted:
  assumes adapted-process M F t_0 Y
 shows adapted-process M F t_0 (\lambda i \xi. X i \xi + Y i \xi)
 using adapted-process.adapted[OF assms] adapted by (unfold-locales) simp
{f lemma} diff-adapted:
  assumes adapted-process M F t_0 Y
 shows adapted-process M F t_0 (\lambda i \xi. X i \xi - Y i \xi)
 using adapted-process.adapted[OF assms] adapted by (unfold-locales) simp
```

```
lemma uminus-adapted: adapted-process MFt_0(-X) using scaleR-right-const-adapted [of
\lambda-. -1] by (simp add: fun-Compl-def)
lemma partial-sum-adapted: adapted-process M F t_0 (\lambda n \xi. \sum i \in \{t_0..n\}. X i \xi)
proof (unfold-locales)
 \mathbf{fix}\ i::\ 'b
 have X j \in borel-measurable (F i) if t_0 \leq j j \leq i for j using that adapted by
  thus (\lambda \xi. \sum i \in \{t_0..i\}. X i \xi) \in borel-measurable (F i) by simp
lemma partial-sum'-adapted: adapted-process M F t_0 (\lambda n \xi. \sum i \in \{t_0... < n\}). X i \xi)
proof (unfold-locales)
 fix i :: 'b
 have X j \in borel-measurable (F i) if t_0 \leq j j < i for j using that adapted by
 thus (\lambda \xi. \sum i \in \{t_0... < i\}. X i \xi) \in borel-measurable (F i) by simp
qed
end
\mathbf{lemma} (in nat\text{-}adapted\text{-}process) partial\text{-}sum\text{-}Suc\text{-}adapted: nat\text{-}adapted\text{-}process M
F(\lambda n \xi. \sum i < n. X(Suc i) \xi)
proof (unfold-locales)
 have X j \in borel-measurable (F i) if j \leq i for j using that adaptedD by blast
 thus (\lambda \xi. \sum i < i. \ X \ (Suc \ i) \ \xi) \in borel-measurable \ (F \ i) by auto
lemma (in filtered-measure) adapted-process-sum:
 assumes \bigwedge i. i \in I \Longrightarrow adapted-process M F t_0 (X i)
 shows adapted-process M F t_0 (\lambda k \xi. \sum i \in I. X i k \xi)
proof -
   fix i k assume i \in I and asm: t_0 \le k
   then interpret adapted-process M F to X i using assms by simp
    have X i k \in borel-measurable M X i k \in borel-measurable (F k) using mea-
surable-from-subalg subalgebras adapted asm by (blast, simp)
 thus ?thesis by (unfold-locales) simp
qed
An adapted process is necessarily a stochastic process.
\mathbf{sublocale}\ adapted-process \subseteq stochastic-process \mathbf{using}\ measurable-from-subalg sub-
algebras adapted by (unfold-locales) blast
sublocale nat-adapted-process \subseteq nat-stochastic-process ..
```

```
\mathbf{sublocale}\ \mathit{real-adapted-process} \subseteq \mathit{real-stochastic-process}\ ..
```

A stochastic process is always adapted to the natural filtration it generates.

```
lemma (in stochastic-process) adapted-process-natural-filtration: adapted-process M (natural-filtration M t_0 X) t_0 X using filtered-measure-natural-filtration by (intro-locales) (auto simp add: natural-filtration-def intro!: adapted-process-axioms.intro measurable-family-vimage-algebra)
```

6.3 Progressively Measurable Process

```
locale progressive-process = filtered-measure M F t_0 for M F t_0 and X :: - \Rightarrow -
\Rightarrow - :: { second-countable-topology, banach} +
 assumes progressive[measurable]: \land t. \ t_0 \leq t \Longrightarrow (\lambda(i, x). \ X \ i \ x) \in borel-measurable
(restrict-space borel \{t_0..t\} \bigotimes_M F t)
begin
lemma progressiveD:
 assumes S \in borel
 shows (\lambda(j, \xi). X j \xi) - 'S \cap (\{t_0..i\} \times space M) \in (restrict\text{-}space borel \{t_0..i\}
 using measurable-sets[OF\ progressive,\ OF\ -\ assms,\ of\ i]
 by (cases t_0 \le i) (auto simp add: space-restrict-space sets-pair-measure space-pair-measure)
end
locale nat-progressive-process = progressive-process M F \theta :: nat X  for M F X
locale real-progressive-process = progressive-process M F \theta :: real X for M F X
{\bf lemma}~({\bf in}~\textit{filtered-measure})~\textit{progressive-process-const-fun}:
 assumes f \in borel-measurable (F t_0)
 shows progressive-process M F t_0 (\lambda -... f)
proof (unfold-locales)
 fix i assume asm: t_0 \leq i
 have f \in borel-measurable (F i) using borel-measurable-mono[OF order.refl asm]
assms by blast
  thus case-prod (\lambda-. f) \in borel-measurable (restrict-space borel \{t_0..i\} \bigotimes_M F(i)
using measurable-compose[OF measurable-snd] by simp
qed
lemma (in filtered-measure) progressive-process-const:
 assumes c \in borel-measurable borel
 shows progressive-process M F t_0 (\lambda i -. c i)
  using assms by (unfold-locales) (auto simp add: measurable-split-conv intro!:
measurable-compose[OF measurable-fst] measurable-restrict-space1)
context progressive-process
```

begin

```
lemma compose-progressive:
    assumes case-prod f \in borel-measurable borel
    shows progressive-process M F t_0 (\lambda i \xi. (f i) (X i \xi))
    fix i assume asm: t_0 \leq i
    have (\lambda(j, \xi), (j, X j \xi)) \in (restrict\text{-space borel } \{t_0, i\} \bigotimes_M F i) \to_M borel \bigotimes_M
borel
         using progressive [OF asm] measurable-fst" [OF measurable-restrict-space1, OF
measurable-id
         by (auto simp add: measurable-pair-iff measurable-split-conv)
    moreover have (\lambda(j, \xi), f_j(X_j \xi)) = case-prod f_j((\lambda(j, y), (j, y))) \circ (\lambda(j, \xi), (j, y)) \circ (\lambda(j, \xi), (j, \xi)) \circ (\lambda(j, \xi),
(j, X j \xi)) by fastforce
    ultimately show (\lambda(j, \xi), (fj), (Xj\xi)) \in borel-measurable (restrict-space borel
\{t_0...i\} \bigotimes_M F(i) using assms by (simp add: borel-prod)
lemma norm-progressive: progressive-process M F t_0 (\lambda i \ \xi. \ norm \ (X \ i \ \xi)) us-
ing measurable-compose [OF progressive borel-measurable-norm] by (unfold-locales)
simp
lemma scaleR-right-progressive:
    assumes progressive-process M F t_0 R
    shows progressive-process M F t_0 (\lambda i \xi. (R i \xi) *_R (X i \xi))
     using progressive-process.progressive[OF assms] by (unfold-locales) (simp add:
progressive assms)
lemma scaleR-right-const-fun-progressive:
    assumes f \in borel-measurable (F t_0)
    shows progressive-process M F t_0 (\lambda i \xi. f \xi *_R (X i \xi))
   using assms by (fast intro: scaleR-right-progressive progressive-process-const-fun)
lemma scaleR-right-const-progressive:
    assumes c \in borel-measurable borel
    shows progressive-process M F t_0 (\lambda i \xi. c i *_R (X i \xi))
   using assms by (fastforce intro: scaleR-right-progressive progressive-process-const)
lemma add-progressive:
     assumes progressive-process M F t_0 Y
    shows progressive-process M F t_0 (\lambda i \xi. X i \xi + Y i \xi)
      using progressive-process.progressive[OF assms] by (unfold-locales) (simp add:
progressive assms)
lemma diff-progressive:
    assumes progressive-process M F t_0 Y
    shows progressive-process M F t_0 (\lambda i \xi. X i \xi - Y i \xi)
     using progressive-process.progressive[OF assms] by (unfold-locales) (simp add:
progressive assms)
```

lemma uminus-progressive: progressive-process MFt_0 (-X) using scaleR-right-const-progressive[of

```
\lambda-. -1] by (simp add: fun-Compl-def)
end
A progressively measurable process is also adapted.
sublocale progressive-process \subseteq adapted-process using measurable-compose-rev[OF]
progressive measurable-Pair1
 unfolding prod.case space-restrict-space
 \mathbf{by} unfold-locales simp
sublocale nat-progressive-process \subseteq nat-adapted-process ...
sublocale real-progressive-process \subseteq real-adapted-process ...
In the discrete setting, adaptedness is equivalent to progressive measurabil-
theorem nat-progressive-iff-adapted: nat-progressive-process MFX \longleftrightarrow nat-adapted-process
MFX
proof (intro iffI)
 assume asm: nat-progressive-process M F X
 interpret nat-progressive-process M F X by (rule asm)
 show nat-adapted-process M F X..
  assume asm: nat-adapted-process M F X
 interpret nat-adapted-process M F X by (rule asm)
 show nat-progressive-process M F X
  proof (unfold-locales, intro borel-measurableI)
   fix S :: 'b \ set \ and \ i :: nat \ assume \ open-S: open \ S
     fix j assume asm: j \leq i
    hence X j - S \cap Space M \in F i using adaptedD[of j, THEN measurable-sets]
space-F open-S by fastforce
      moreover have case-prod X - 'S \cap \{j\} \times space M = \{j\} \times (X j - `S \cap S)
space M) for j by fast
     moreover have \{j :: nat\} \in restrict\text{-}space borel <math>\{0..i\} using asm by (simp)
add: sets-restrict-space-iff)
      ultimately have case-prod X - 'S \cap \{j\} \times space M \in restrict\text{-space borel}
\{\theta..i\} \bigotimes_M F i  by simp
   }
   hence (\lambda j. \ (\lambda(x, y). \ X \ x \ y) - `S \cap \{j\} \times space \ M) \ `\{..i\} \subseteq restrict\text{-space borel}
\{0..i\} \bigotimes_M F i  by blast
   moreover have case-prod X - 'S \cap space (restrict-space borel \{0..i\} \bigotimes_M F
i) = (\bigcup j \le i. \ case-prod \ X - S \cap \{j\} \times space \ M) \ unfolding \ space-pair-measure
space-restrict-space space-F by force
    ultimately show case-prod X - 'S \cap space (restrict-space borel \{0...i\} \bigotimes_{M}
F(i) \in restrict\text{-space borel } \{0..i\} \bigotimes_{M} F(i) \text{ by } (metis\ sets.countable-UN)
 qed
\mathbf{qed}
```

6.4 Predictable Process

```
We introduce the constant \Sigma_P to denote the predictable sigma algebra.
{\bf context}\ {\it linearly-filtered-measure}
begin
definition \Sigma_P :: (b \times a) measure where predictable-sigma: \Sigma_P \equiv sigma \ (\{t_0..\}\}
\times \ space \ M) \ (\{\{s<..t\} \times A \mid A \ s \ t. \ A \in F \ s \land t_0 \le s \land s < t\} \cup \{\{t_0\} \times A \mid A.
A \in F t_0
lemma space-predictable-sigma[simp]: space \Sigma_P = (\{t_0..\} \times space\ M) unfolding
predictable-sigma space-measure-of-conv by blast
lemma sets-predictable-sigma: sets \Sigma_P = sigma-sets (\{t_0..\} \times space\ M) (\{\{s<..t\}\}
\times A \mid A \ s \ t. \ A \in F \ s \land t_0 \le s \land s < t \} \cup \{ \{t_0\} \times A \mid A. \ A \in F \ t_0 \} )
 unfolding predictable-sigma using space-F sets.sets-into-space by (subst sets-measure-of)
fastforce+
{\bf lemma}\ measurable\text{-}predictable\text{-}sigma\text{-}snd\text{:}
  assumes countable \mathcal{I} \mathcal{I} \subseteq \{\{s<..t\} \mid s \ t. \ t_0 \leq s \land s < t\} \ \{t_0<..\} \subseteq (\bigcup \mathcal{I})
  shows snd \in \Sigma_P \to_M F t_0
proof (intro measurableI)
  fix S :: 'a set assume asm: S \in F t_0
  have countable: countable ((\lambda I.\ I \times S)\ '\mathcal{I}) using assms(1) by blast
  have (\lambda I. \ I \times S) '\mathcal{I} \subseteq \{\{s < ...t\} \times A \mid A \ s \ t. \ A \in F \ s \land t_0 \leq s \land s < t\} using
sets-F-mono[OF \ order-refl, \ THEN \ subsetD, \ OF - \ asm] \ assms(2) \ \mathbf{by} \ blast
 hence (\bigcup I \in \mathcal{I}. \ I \times S) \cup \{t_0\} \times S \in \Sigma_P  unfolding sets-predictable-sigma using
asm\ \mathbf{by}\ (intro\ sigma-sets-Un[OF\ sigma-sets-UNION[OF\ countable]\ sigma-sets.Basic]
sigma-sets.Basic) blast+
 moreover have snd - S \cap space \Sigma_P = \{t_0..\} \times S \text{ using } sets.sets-into-space [OF] \}
asm] by fastforce
  moreover have \{t_0\} \cup \{t_0 < ...\} = \{t_0...\} by auto
  moreover have (\bigcup I \in \mathcal{I}. \ I \times S) \cup \{t_0\} \times S = \{t_0..\} \times S \text{ using } assms(2,3)
calculation(3) by fastforce
  ultimately show snd - Snace \Sigma_P \in \Sigma_P by argo
qed (auto)
lemma measurable-predictable-sigma-fst:
  assumes countable \mathcal{I} \mathcal{I} \subseteq \{\{s<..t\} \mid s \ t. \ t_0 \leq s \land s < t\} \ \{t_0<..\} \subseteq (\bigcup \mathcal{I})
  shows fst \in \Sigma_P \to_M borel
proof -
  have A \times space \ M \in sets \ \Sigma_P \ \text{if} \ A \in sigma-sets \ \{t_0..\} \ \{\{s<..t\} \mid s \ t. \ t_0 \leq s \land s \}
\langle t \rangle for A unfolding sets-predictable-sigma using that
  proof (induction rule: sigma-sets.induct)
    case (Basic\ a)
    thus ?case using space-F sets.top by blast
  next
    case (Compl\ a)
    have (\{t_0..\} - a) \times space M = \{t_0..\} \times space M - a \times space M by blast
```

```
then show ?case using Compl(2)[THEN sigma-sets.Compl] by presburger
  next
   case (Union \ a)
   have [\ ] (range a) \times space M = [\ ] (range (\lambda i.\ a\ i \times space\ M)) by blast
   then show ?case using Union(2)[THEN sigma-sets.Union] by presburger
  moreover have restrict-space borel \{t_0..\} = sigma \{t_0..\} \{\{s < ..t\} \mid s \ t. \ t_0 \le s
\land s < t
 proof -
    have sigma-sets \{t_0..\} ((\cap)\ \{t_0..\}\ 'sigma-sets\ UNIV\ (range\ greaterThan)) =
sigma-sets \{t_0..\} \{\{s<..t\} | s \ t. \ t_0 \le s \land s < t\}
   proof (intro sigma-sets-eqI; clarify)
     fix A :: 'b \text{ set assume } asm: A \in sigma-sets UNIV (range greaterThan)
     thus \{t_0..\} \cap A \in sigma\text{-sets } \{t_0..\} \{\{s<..t\} \mid s \ t. \ t_0 \leq s \land s < t\}
     proof (induction rule: sigma-sets.induct)
       case (Basic\ a)
       then obtain s where s: a = \{s < ...\} by blast
       show ?case
       proof (cases t_0 \leq s)
         case True
         hence *: \{t_0..\} \cap a = (\bigcup i \in \mathcal{I}. \{s<..\} \cap i) using s \ assms(3) by force
         have ((\cap) \{s<...\} `\mathcal{I}) \subseteq sigma-sets \{t_0...\} \{\{s<...t\} \mid s \ t. \ t_0 \leq s \land s < t\}
         proof (clarify)
           fix A assume A \in \mathcal{I}
         then obtain s' t' where A: A = \{s' < ...t'\}\ t_0 \le s' s' < t' using assms(2)
by blast
           hence \{s<...\} \cap A = \{max \ s \ s'<..t'\} by fastforce
           moreover have t_0 \leq max \ s \ ' using A True by linarith
           moreover have \max s \ s' < t' \ \text{if} \ s < t' \ \text{using} \ A \ that \ \text{by} \ linarith
           moreover have \{s<..\} \cap A = \{\} if \neg s < t' using A that by force
           ultimately show \{s<..\} \cap A \in sigma\text{-}sets \{t_0..\} \{\{s<..t\} \mid s \ t. \ t_0 \leq s
\land s < t} by (cases s < t') (blast, simp add: sigma-sets. Empty)
         thus ?thesis unfolding * using assms(1) by (intro sigma-sets-UNION)
auto
       next
         case False
         hence \{t_0..\} \cap a = \{t_0..\} using s by force
         thus ?thesis using sigma-sets-top by auto
       qed
     next
       case (Compl\ a)
       have \{t_0..\} \cap (UNIV - a) = \{t_0..\} - (\{t_0..\} \cap a) by blast
       then show ?case using Compl(2)[THEN sigma-sets.Compl] by presburger
     next
       case (Union \ a)
       have \{t_0..\} \cap \bigcup (range \ a) = \bigcup (range \ (\lambda i. \{t_0..\} \cap a \ i)) by blast
       then show ?case using Union(2)[THEN sigma-sets.Union] by presburger
     qed (simp add: sigma-sets.Empty)
```

```
next
     fix s \ t assume asm: t_0 \le s \ s < t
     hence *: \{s<...t\} = \{s<...\} \cap (\{t_0...\} - \{t<...\}) by force
    have \{s<...\} \in sigma-sets \{t_0...\} ((\cap) \{t_0...\} 'sigma-sets UNIV (range greaterThan))
using asm by (intro sigma-sets.Basic) auto
     moreover have \{t_0..\} - \{t<..\} \in sigma-sets \{t_0..\} ((\cap) \{t_0..\} `sigma-sets
UNIV (range greaterThan)) using asm by (intro sigma-sets.Compl sigma-sets.Basic)
     ultimately show \{s<..t\} \in sigma-sets \{t_0..\} ((\cap) \{t_0..\} ' sigma-sets UNIV
(range\ greaterThan))\ \mathbf{unfolding}*Int-range-binary[of\ \{s<..\}]\ \mathbf{by}\ (intro\ sigma-sets-Inter[OF\ sets])
- binary-in-sigma-sets]) auto
    thus ?thesis unfolding borel-Ioi restrict-space-def emeasure-sigma by (force
intro: sigma-eqI)
  qed
 ultimately have restrict-space borel \{t_0..\} \bigotimes_M sigma (space M) \{\} \subseteq sets \Sigma_P
   unfolding sets-pair-measure space-restrict-space space-measure-of-conv
   using space-predictable-sigma sets.sigma-algebra-axioms[of \Sigma_P]
   by (intro sigma-algebra.sigma-sets-subset) (auto simp add: sigma-sets-empty-eq
sets-measure-of-conv)
  moreover have space (restrict-space borel \{t_0..\} \bigotimes_M sigma (space M) \{\}) =
space \Sigma_P by (simp add: space-pair-measure)
  moreover have fst \in restrict\text{-}space \ borel\ \{t_0..\}\ \bigotimes_{M} \ sigma\ (space\ M)\ \{\} \rightarrow_{M}
borel by (fastforce intro: measurable-fst" [OF measurable-restrict-space1, of \lambda x. x])
 ultimately show ?thesis by (meson borel-measurable-subalgebra)
qed
end
locale predictable-process = linearly-filtered-measure\ M\ F\ t_0 for M\ F\ t_0 and X:
- \Rightarrow - \Rightarrow - :: \{second\text{-}countable\text{-}topology, banach} +
 assumes predictable: (\lambda(t, x). X t x) \in borel-measurable \Sigma_P
begin
{f lemmas}\ predictable D = measurable-sets[OF\ predictable,\ unfolded\ space-predictable-sigma]
end
locale nat-predictable-process = predictable-process M F 0 :: nat X for M F X
locale real-predictable-process = predictable-process M F 0 :: real X \text{ for } M F X
lemma (in nat-filtered-measure) measurable-predictable-sigma-snd':
 shows snd \in \Sigma_P \to_M F \theta
 by (intro measurable-predictable-sigma-snd[of range (\lambda x. {Suc x})]) (force | simp
add: greaterThan-\theta)+
```

```
lemma (in nat-filtered-measure) measurable-predictable-sigma-fst':
 shows fst \in \Sigma_P \to_M borel
 by (intro measurable-predictable-sigma-fst[of range (\lambda x. \{Suc\ x\})]) (force | simp
add: qreaterThan-\theta)+
lemma (in real-filtered-measure) measurable-predictable-sigma-snd':
 shows snd \in \Sigma_P \to_M F \theta
 using real-arch-simple by (intro measurable-predictable-sigma-snd) of range (\lambda x::nat.
\{0 < ... real (Suc x)\}\} (fastforce intro: add-increasing)+
lemma (in real-filtered-measure) measurable-predictable-sigma-fst':
 shows fst \in \Sigma_P \to_M borel
 using real-arch-simple by (intro measurable-predictable-sigma-fst of range (\lambda x::nat.
\{0 < ... real (Suc x)\}\} (fastforce intro: add-increasing)+
lemma (in linearly-filtered-measure) predictable-process-const-fun:
 assumes snd \in \Sigma_P \to_M F t_0 f \in borel\text{-}measurable (F t_0)
   shows predictable-process M F t_0 (\lambda -... f)
  using measurable-compose-rev[OF\ assms(2)]\ assms(1) by (unfold-locales) (auto
simp add: measurable-split-conv)
lemma (in nat-filtered-measure) predictable-process-const-fun'[intro]:
 assumes f \in borel-measurable (F \ \theta)
 shows nat-predictable-process M F (\lambda-. f)
 using assms by (intro predictable-process-const-fun OF measurable-predictable-sigma-snd',
THEN nat-predictable-process.intro])
lemma (in real-filtered-measure) predictable-process-const-fun'[intro]:
  assumes f \in borel-measurable (F \ 0)
 shows real-predictable-process M F (\lambda - f)
 using assms by (intro predictable-process-const-fun OF measurable-predictable-sigma-snd',
THEN real-predictable-process.intro])
lemma (in linearly-filtered-measure) predictable-process-const:
  assumes fst \in borel-measurable \Sigma_P c \in borel-measurable borel
 shows predictable-process M F t_0 (\lambda i -. c i)
  using assms by (unfold-locales) (simp add: measurable-split-conv)
lemma (in linearly-filtered-measure) predictable-process-const-const[intro]:
  shows predictable-process M F t_0 (\lambda - - c)
 by (unfold-locales) simp
lemma (in nat-filtered-measure) predictable-process-const'[intro]:
  assumes c \in borel-measurable borel
 shows nat-predictable-process M F (\lambda i -. c i)
 using assms by (intro predictable-process-const[OF measurable-predictable-sigma-fst',
```

```
THEN nat-predictable-process.intro])
lemma (in real-filtered-measure) predictable-process-const'[intro]:
 assumes c \in borel-measurable borel
 shows real-predictable-process M F (\lambda i - c i)
 using assms by (intro predictable-process-const[OF measurable-predictable-siqma-fst',
THEN real-predictable-process.intro])
context predictable-process
begin
lemma compose-predictable:
 assumes fst \in borel-measurable \Sigma_P case-prod f \in borel-measurable borel
 shows predictable-process M F t_0 (\lambda i \xi. (f i) (X i \xi))
proof
 have (\lambda(i, \xi), (i, X i \xi)) \in \Sigma_P \to_M borel \bigotimes_M borel using predictable assms(1)
by (auto simp add: measurable-pair-iff measurable-split-conv)
  moreover have (\lambda(i, \xi), f(i, X(i, \xi))) = case-prod f(i, \xi), (i, X(i, \xi)) by
 ultimately show (\lambda(i, \xi), f(X i \xi)) \in borel-measurable \Sigma_P unfolding borel-prod
using assms by simp
qed
lemma norm-predictable: predictable-process M F t_0 (\lambda i \xi. norm (X i \xi)) using
measurable-compose[OF predictable borel-measurable-norm]
 by (unfold-locales) (simp add: prod.case-distrib)
\mathbf{lemma}\ scaleR	ext{-}right	ext{-}predictable:
 assumes predictable-process M F t_0 R
 shows predictable-process M F t_0 (\lambda i \xi. (R i \xi) *_R (X i \xi))
  using predictable predictable-process.predictable[OF assms] by (unfold-locales)
(auto simp add: measurable-split-conv)
\mathbf{lemma}\ scaleR\text{-}right\text{-}const\text{-}fun\text{-}predictable\text{:}
 assumes snd \in \Sigma_P \to_M F t_0 f \in borel\text{-}measurable (F t_0)
 shows predictable-process M F t_0 (\lambda i \xi. f \xi *_R (X i \xi))
 using assms by (fast intro: scaleR-right-predictable predictable-process-const-fun)
lemma scaleR-right-const-predictable:
 assumes fst \in borel-measurable \Sigma_P c \in borel-measurable borel
 shows predictable-process M F t_0 (\lambda i \xi. c i *_R (X i \xi))
 using assms by (fastforce intro: scaleR-right-predictable predictable-process-const)
lemma scaleR-right-const'-predictable: predictable-process M F t_0 (\lambda i \xi. c *_R (X i)
 by (fastforce intro: scaleR-right-predictable)
lemma add-predictable:
```

assumes predictable-process $M F t_0 Y$

```
shows predictable-process M F t_0 (\lambda i \xi. X i \xi + Y i \xi)
  using predictable predictable-process.predictable[OF assms] by (unfold-locales)
(auto simp add: measurable-split-conv)
lemma diff-predictable:
  assumes predictable-process M F t_0 Y
  shows predictable-process M F t_0 (\lambda i \xi. X i \xi - Y i \xi)
  using predictable predictable-process.predictable[OF assms] by (unfold-locales)
(auto simp add: measurable-split-conv)
lemma uminus-predictable: predictable-process MF t_0 (-X) using scaleR-right-const'-predictable[of
-1] by (simp add: fun-Compl-def)
end
Every predictable process is also progressively measurable.
sublocale predictable-process \subseteq progressive-process
proof (unfold-locales)
  fix i :: 'b assume asm: t_0 \leq i
  {
   fix S :: (b \times a) set assume S \in \{\{s < ... t\} \times A \mid A \mid t \in F \mid s \land t_0 \leq s \land s \}
\{t\} \cup \{\{t_0\} \times A \mid A. A \in F \ t_0\}
   hence (\lambda x.\ x) - 'S \cap (\{t_0..i\} \times space\ M) \in restrict\text{-}space\ borel\ \{t_0..i\} \bigotimes_M F
i
      assume S \in \{\{s < ...t\} \times A \mid A \ s \ t. \ A \in F \ s \land t_0 \le s \land s < t\}
      then obtain s \ t \ A where S-is: S = \{s < ...t\} \times A \ t_0 \le s \ s < t \ A \in F \ s by
blast
       hence (\lambda x. \ x) - S \cap (\{t_0...i\} \times space \ M) = \{s < ... min \ i \ t\} \times A \ using
sets.sets-into-space[OF\ S-is(4)] by auto
     then show ?thesis using S-is sets-F-mono[of s i] by (cases s \leq i) (fastforce
simp add: sets-restrict-space-iff)+
   next
      assume S \in \{\{t_0\} \times A \mid A. A \in F t_0\}
      then obtain A where S-is: S = \{t_0\} \times A A \in F t_0 \text{ by } blast
    hence (\lambda x.\ x) - S \cap (\{t_0..i\} \times space\ M) = \{t_0\} \times A \text{ using } asm\ sets.sets-into-space}[OF]
S-is(2)] by auto
     thus ?thesis using S-is(2) sets-F-mono[OF order-reft asm] asm by (fastforce
simp add: sets-restrict-space-iff)
   qed
  hence (\lambda x. x) - S \cap space (restrict-space borel \{t_0...i\} \bigotimes_M F i) \in restrict-space
borel \{t_0..i\} \bigotimes_M F i by (simp \ add: space-pair-measure \ space-F[OF \ asm])
 moreover have \{\{s < ...t\} \times A \mid A \text{ s. } t. A \in sets (F \text{ s}) \land t_0 \leq s \land s < t\} \cup \{\{t_0\}\}\}
\times A \mid A. A \in sets(F \mid t_0) \subseteq Pow(\{t_0...\} \times spaceM) using sets.sets-into-space by
force
  ultimately have (\lambda x. \ x) \in restrict\text{-space borel } \{t_0..i\} \bigotimes_M F \ i \to_M \Sigma_P \text{ us-}
ing space-F[OF asm] by (intro\ measurable-sigma-sets[OF\ sets-predictable-sigma])
```

(fast, force simp add: space-pair-measure)

```
thus case-prod X \in borel-measurable (restrict-space borel \{t_0...i\} \bigotimes_M F(i)) using
predictable by simp
qed
sublocale nat-predictable-process \subseteq nat-progressive-process ...
sublocale real-predictable-process \subseteq real-progressive-process ...
The following lemma characterizes predictability in a discrete-time setting.
lemma (in nat-filtered-measure) sets-in-filtration:
 assumes (\bigcup i. \{i\} \times A \ i) \in \Sigma_P
 shows A (Suc i) \in F i A 0 \in F 0
  using assms unfolding sets-predictable-sigma
proof (induction (\bigcup i. \{i\} \times A \ i) arbitrary: A)
 case Basic
   assume \exists S. (\bigcup i. \{i\} \times A \ i) = \{0\} \times S
   then obtain S where S: (\bigcup i. \{i\} \times A \ i) = \{bot\} \times S \ unfolding \ bot-nat-def
by blast
   hence S \in F bot using Basic by (fastforce simp add: times-eq-iff bot-nat-def)
   moreover have A \ i = \{\} if i \neq bot for i using that S by blast
   moreover have A \ bot = S \ using \ S \ by \ blast
   ultimately have A (Suc i) \in F i A \theta \in F \theta for i unfolding bot-nat-def by
(auto simp add: bot-nat-def)
 note * = this
   assume \nexists S. (\bigcup i. \{i\} \times A \ i) = \{0\} \times S
   then obtain s t B where B: (\bigcup i. \{i\} \times A \ i) = \{s < ... t\} \times B \ B \in sets \ (F \ s)
s < t using Basic by auto
   hence A \ i = B \ \text{if} \ i \in \{s < ... t\} \ \text{for} \ i \ \text{using} \ that \ \text{by} \ fast
   moreover have A i = \{\} if i \notin \{s < ... t\} for i using B that by fastforce
   ultimately have A (Suc i) \in F i A \theta \in F \theta for i unfolding bot-nat-def using
B sets-F-mono by (auto simp add: bot-nat-def) (metis less-Suc-eq-le sets.empty-sets
subset-eq)
 }
 note ** = this
 show A (Suc i) \in sets (F i) A \theta \in sets (F \theta) using *(1)[of i] *(2) **(1)[of i]
**(2) by blast+
\mathbf{next}
 case Empty
   case 1
   then show ?case using Empty by simp
 next
   case 2
   then show ?case using Empty by simp
  }
next
 case (Compl\ a)
```

```
\mathbf{have}\ a\text{-}in:\ a\subseteq\{0..\}\times space\ M\ \mathbf{using}\ Compl(1)\ sets.sets\text{-}into\text{-}space\ sets\text{-}predictable\text{-}sigma
space-predictable-sigma by metis
  hence A-in: A i \subseteq space\ M for i using Compl(4) by blast
 have a: a = \{0..\} \times space \ M - (\bigcup i. \{i\} \times A \ i)  using a-in Compl(4) by blast
  also have ... = -(\bigcap j. - (\{j\} \times (space M - A j))) by blast
  also have ... = (\bigcup j. \{j\} \times (space M - A j)) by blast
  finally have *: (space\ M-A\ (Suc\ i))\in F\ i\ (space\ M-A\ 0)\in F\ 0 using
Compl(2,3) by auto
  {
   case 1
     then show ?case using * A-in by (metis bot-nat-0.extremum double-diff
sets.Diff sets.top sets-F-mono sets-le-imp-space-le space-F)
 next
   case 2
     then show ?case using * A-in by (metis bot-nat-0.extremum double-diff
sets.Diff sets.top sets-F-mono sets-le-imp-space-le space-F)
  }
next
  case (Union \ a)
  have a-in: a \in \{0...\} \times space \ M \ \text{for} \ i \ \text{using} \ Union(1) \ sets.sets-into-space
sets-predictable-sigma space-predictable-sigma by metis
  hence A-in: A i \subseteq space \ M for i \ using \ Union(4) by blast
  have snd \ x \in snd \ (a \ i \cap (\{fst \ x\} \times space \ M)) \ \textbf{if} \ x \in a \ i \ \textbf{for} \ i \ x \ \textbf{using} \ that
a-in by fastforce
  hence a-i: a i = (\bigcup j. \{j\} \times (snd \cdot (a \ i \cap (\{j\} \times space \ M)))) for i by force
  have A-i: A i = snd ' (\bigcup (range a) \cap (\{i\} \times space M)) for i unfolding
Union(4) using A-in by force
 have *: snd '(a \ j \cap (\{Suc\ i\} \times space\ M)) \in F\ i\ snd '(a \ j \cap (\{\emptyset\} \times space\ M))
\in F \ 0 \ \text{for} \ j \ \text{using} \ Union(2,3)[OF \ a-i] \ \text{by} \ auto
  {
   case 1
   have (\bigcup j. \ snd \ (a \ j \cap (\{Suc \ i\} \times space \ M))) \in F \ i \ using * by \ fast
    moreover have (\bigcup j. \ snd \ (a \ j \cap (\{Suc \ i\} \times space \ M))) = snd \ (\bigcup \ (range)
a) \cap (\{Suc\ i\} \times space\ M)) by fast
   ultimately show ?case using A-i by metis
  next
   case 2
   have (\bigcup j. \ snd \ (a \ j \cap (\{0\} \times space \ M))) \in F \ 0 \ using * by fast
   moreover have (\bigcup j. \ snd \ (a \ j \cap (\{0\} \times space \ M))) = snd \ (\bigcup \ (range \ a) \cap \{0\} \times space \ M))
(\{\theta\} \times space\ M)) by fast
   ultimately show ?case using A-i by metis
  }
qed
This leads to the following useful fact.
lemma (in nat-predictable-process) adapted-Suc: nat-adapted-process M F (\lambda i. X)
(Suc\ i)
proof (unfold-locales, intro borel-measurableI)
  fix S :: 'b \ set \ and \ i \ assume \ open-S: \ open \ S
```

```
have \{Suc\ i\} = \{i < ... Suc\ i\} by fastforce
    hence \{Suc\ i\} \times space\ M \in \Sigma_P \text{ using } space\text{-}F[symmetric,\ of\ i] \text{ unfolding}
sets-predictable-sigma by (intro sigma-sets.Basic) blast
    moreover have case-prod X - S \cap (UNIV \times space M) \in \Sigma_P unfolding
atLeast-0[symmetric] using open-S by (intro predictableD, simp add: borel-open)
    ultimately have case-prod X - S \cap (\{Suc\ i\} \times space\ M) \in \Sigma_P unfolding
sets-predictable-sigma using space-F sets.sets-into-space
          by (subst Times-Int-distrib1 of {Suc i} UNIV space M, simplified], subst
inf.commute, subst Int-assoc[symmetric], subst Int-range-binary)
            (intro\ sigma-sets-Inter\ binary-in-sigma-sets,\ fast)+
  moreover have case-prod X - S \cap (\{Suc\ i\} \times space\ M) = \{Suc\ i\} \times (X (Suc\ i) + Suc\ i) \times (X (Suc
i) - 'S \cap space M) by (auto simp add: le-Suc-eq)
  moreover have ... = (\bigcup j. \{j\} \times (if j = Suc \ i \ then \ (X \ (Suc \ i) - `S \cap space \ M)
else {})) by (force split: if-splits)
  ultimately have (\bigcup j. \{j\} \times (if j = Suc \ i \ then \ (X \ (Suc \ i) - `S \cap space \ M) \ else
\{\})) \in \Sigma_P by argo
    thus X (Suc i) - 'S \cap space (F i) \in sets (F i) using sets-in-filtration[of \lambda j.
if j = Suc \ i \ then \ (X \ (Suc \ i) - `S \cap space \ M) \ else \ \{\}] \ space-F[OF \ zero-le] \ by
presburger
qed
\textbf{theorem} \ \textit{nat-predictable-process-iff: nat-predictable-process} \ \textit{MFX} \longleftrightarrow \textit{nat-adapted-process}
M F (\lambda i. X (Suc i)) \wedge X \theta \in borel\text{-}measurable (F \theta)
proof (intro iffI)
   assume asm: nat-adapted-process M F (\lambda i.~X~(Suc~i)) \wedge~X~\theta \in borel-measurable
(F \theta)
   interpret nat-adapted-process M F \lambda i. X (Suc i) using asm by blast
   have (\lambda(x, y). X x y) \in borel\text{-}measurable \Sigma_P
   proof (intro borel-measurableI)
       fix S :: 'b \ set \ assume \ open-S: \ open \ S
       have \{i\} \times (X \ i - S \cap space \ M) \in sets \ \Sigma_P \ for \ i
       proof (cases i)
          case \theta
          then show ?thesis unfolding sets-predictable-sigma
            using measurable-sets[OF - borel-open[OF open-S], of X 0 F 0] asm by auto
       next
          case (Suc i)
          have \{Suc\ i\} = \{i < ... Suc\ i\} by fastforce
          then show ?thesis unfolding sets-predictable-sigma
              using measurable-sets[OF adapted borel-open[OF open-S], of i]
              by (intro sigma-sets.Basic, auto simp add: Suc)
      moreover have (\lambda(x, y). \ X \ x \ y) - `S \cap space \ \Sigma_P = (\bigcup i. \ \{i\} \times (X \ i - `S \cap S))
space M)) by fastforce
       ultimately show (\lambda(x, y). X x y) - S \cap space \Sigma_P \in sets \Sigma_P by simp
    thus nat-predictable-process M F X by (unfold-locales)
next
   assume asm: nat-predictable-process M F X
```

```
interpret nat-predictable-process M F X by (rule\ asm) show nat-adapted-process M F (\lambda i.\ X\ (Suc\ i)) \land X\ \theta \in borel-measurable\ (F\ \theta) using adapted\text{-}Suc\ by\ simp \mathbf{qed}
```

end

theory Martingale imports Stochastic-Process Conditional-Expectation-Banach begin

7 Martingales

The following locales are necessary for defining martingales.

7.1 Additional Locale Definitions

 $\begin{array}{l} \textbf{locale} \ sigma-finite-adapted-process = sigma-finite-filtered-measure} \ M \ F \ t_0 \ X \ \textbf{for} \ M \ F \ t_0 \ X \\ \end{array}$

 $\label{locale} \begin{subarray}{l} \textbf{locale} \ nat\text{-}sigma\text{-}finite\text{-}adapted\text{-}process} \ M\ F\ 0 :: nat \\ X\ \textbf{for}\ M\ F\ X \end{subarray}$

 $\label{locale} \textbf{locale} \ \textit{real-sigma-finite-adapted-process} \ \textit{M} \ \textit{F} \ \textit{0} \ :: \\ \textit{real} \ \textit{X} \ \textbf{for} \ \textit{M} \ \textit{F} \ \textit{X}$

sublocale nat-sigma-finite-adapted-process \subseteq nat-sigma-finite-filtered-measure .. sublocale real-sigma-finite-adapted-process \subseteq real-sigma-finite-filtered-measure ..

locale finite-adapted-process = finite-filtered-measure $M F t_0 + adapted$ -process $M F t_0 X$ for $M F t_0 X$

 $\mathbf{sublocale}\ finite-adapted-process\subseteq sigma-finite-adapted-process$..

 $\label{eq:locale_note} \textbf{locale} \ \textit{nat-finite-adapted-process} \ \textit{M} \ \textit{F} \ \textit{0} \ :: \ \textit{nat} \ \textit{X} \ \textbf{for} \ \textit{M} \ \textit{F} \ \textit{X}$

 $\label{locale} \textbf{locale} \ \textit{real-finite-adapted-process} \ \textit{M} \ \textit{F} \ \textit{0} \ :: \ \textit{real} \ \textit{X} \ \textbf{for} \ \textit{M} \\ \textit{F} \ \textit{X}$

sublocale nat-finite-adapted-process \subseteq nat-sigma-finite-adapted-process .. sublocale real-finite-adapted-process \subseteq real-sigma-finite-adapted-process ..

locale sigma-finite-adapted-process-order = sigma-finite-adapted-process $M F t_0 X$ for $M F t_0$ and $X :: - \Rightarrow - \Rightarrow - :: \{order-topology, ordered-real-vector\}$

 $\label{locale-notes} \begin{subarray}{l} \textbf{locale} & \textit{nat-sigma-finite-adapted-process-order} \\ \textit{M} \ \textit{F} \ \textit{0} :: \textit{nat} \ \textit{X} \ \textbf{for} \ \textit{M} \ \textit{F} \ \textit{X} \end{subarray}$

 $\label{locale} \begin{subarray}{l} \textbf{locale} \end{subarray} real-sigma-finite-adapted-process-order = sigma-finite-adapted-process-order \\ M \end{subarray} \begin{subarray}{l} A \end{subarray} \begin{subarray}{l} \textbf{locale} \end{subarray} real X \end{subarray} \begin{subarray}{l} \textbf{locale} \end{subarray} real-sigma-finite-adapted-process-order \\ \textbf{locale} \end{subarray} \begin{subarray}{l} \textbf{locale} \end{subarray} \begin{suba$

sublocale nat-sigma-finite-adapted-process-order \subseteq nat-sigma-finite-adapted-process ...

 $\textbf{sublocale}\ real\text{-}sigma\text{-}finite\text{-}adapted\text{-}process\text{-}order \subseteq real\text{-}sigma\text{-}finite\text{-}adapted\text{-}process$..

locale finite-adapted-process-order = finite-adapted-process $M F t_0 X$ for $M F t_0$ and $X :: - \Rightarrow - \Rightarrow - :: \{order-topology, ordered-real-vector\}$

 $\label{locale} \begin{subarray}{l} \textbf{locale} \ nat\text{-}finite\text{-}adapted\text{-}process\text{-}order = finite\text{-}adapted\text{-}process\text{-}order M F 0 :: nat } X \ \textbf{for} \ M F X \end{subarray}$

 $\label{locale} \textbf{locale} \ \textit{real-finite-adapted-process-order} \ \textit{ } F \ \textit{0} \ :: \ \textit{real} \ \textit{X} \ \textbf{for} \ \textit{M} \ \textit{F} \ \textit{X}$

 $\begin{array}{l} \textbf{sublocale} \ nat\text{-}finite\text{-}adapted\text{-}process\text{-}order \subseteq nat\text{-}sigma\text{-}finite\text{-}adapted\text{-}process\text{-}order \\ .. \\ \textbf{sublocale} \ real\text{-}finite\text{-}adapted\text{-}process\text{-}order \subseteq real\text{-}sigma\text{-}finite\text{-}adapted\text{-}process\text{-}order \\ \end{array}$

locale $sigma-finite-adapted-process-linorder = sigma-finite-adapted-process-order <math>M \ F \ t_0 \ X \ for \ M \ F \ t_0 \ and \ X :: - \Rightarrow - \Rightarrow - :: \{linorder-topology\}$

 $\label{locale} \textbf{locale} \ \textit{nat-sigma-finite-adapted-process-linorder} = \textit{sigma-finite-adapted-process-linorder} \\ \textit{M} \ \textit{F} \ \textit{0} :: \textit{nat} \ \textit{X} \ \textbf{for} \ \textit{M} \ \textit{F} \ \textit{X}$

 $\label{locale} \textbf{locale} \ \textit{real-sigma-finite-adapted-process-linorder} = \textit{sigma-finite-adapted-process-linorder} \\ M \ F \ 0 :: \textit{real} \ X \ \textbf{for} \ M \ F \ X$

locale finite-adapted-process-linorder = finite-adapted-process-order $M F t_0 X$ for $M F t_0$ and $X :: - \Rightarrow - \Rightarrow - :: \{linorder-topology\}$

 $\label{locale} \begin{subarray}{l} \textbf{locale} \end{subarray} \begin{subarray}{l} \textbf{locale} \end{subarray} \begin{subarray}{l} \textbf{nat-finite-adapted-process-linorder} \end{subarray} \begin{subarray}{l} \textbf{for} \end{subarray} \begin{subarray}{l} \textbf{F} \end{subarray} \begin{subarray}{l} \textbf{for} \end{subarray} \begin{subarray}{l} \textbf{for$

 $\begin{array}{l} \textbf{locale} \ \ real\mbox{-}finite\mbox{-} adapted\mbox{-}process\mbox{-}linorder \ = \mbox{-}finite\mbox{-} adapted\mbox{-}process\mbox{-}linorder \ M\ F\ 0 \\ :: \ real\ X\ \ \textbf{for}\ \ M\ F\ X \\ \end{array}$

 ••

7.2 Martingale

locale martingale = sigma-finite-adapted-process +

```
assumes integrable: \bigwedge i. t_0 \leq i \Longrightarrow integrable \ M \ (X \ i)
     and martingale-property: \bigwedge i j. t_0 \leq i \Longrightarrow i \leq j \Longrightarrow AE \xi in M. X i \xi =
cond\text{-}exp\ M\ (F\ i)\ (X\ j)\ \xi
locale martingale-order = martingale M F t_0 X for M F t_0 and X :: - \Rightarrow -
:: \{order-topology, ordered-real-vector\}
locale martingale-linorder = martingale M F t_0 X for M F t_0 and X :: - \Rightarrow - \Rightarrow
- :: {linorder-topology, ordered-real-vector}
sublocale martingale-linorder \subseteq martingale-order ..
lemma (in sigma-finite-filtered-measure) martingale-const-fun[intro]:
  assumes integrable M f f \in borel-measurable (F t_0)
 shows martingale M F t_0 (\lambda-. f)
 using assms sigma-finite-subalgebra.cond-exp-F-meas[OF - assms(1), THEN AE-symmetric]
borel-measurable-mono
 by (unfold-locales) blast+
lemma (in sigma-finite-filtered-measure) martingale-cond-exp[intro]:
 assumes integrable M f
 shows martingale M F t_0 (\lambda i. cond\text{-}exp M (F i) f)
 {f using}\ sigma-finite-subalgebra. borel-measurable-cond-exp'\ borel-measurable-cond-exp
 by (unfold-locales) (auto intro: sigma-finite-subalgebra.cond-exp-nested-subalg[OF]
- assms] simp add: subalgebra-F subalgebras)
corollary (in sigma-finite-filtered-measure) martingale-zero[intro]: martingale M
F t_0 (\lambda - - \cdot \cdot \theta) by fastforce
corollary (in finite-filtered-measure) martingale-const[intro]: martingale M F t_0
(\lambda- -. c) by fastforce
7.3
       Submartingale
locale\ submartingale = sigma-finite-adapted-process-order +
 assumes integrable: \bigwedge i. t_0 \leq i \Longrightarrow integrable \ M \ (X \ i)
     cond\text{-}exp\ M\ (F\ i)\ (X\ j)\ \xi
locale submartingale-linorder = submartingale\ M\ F\ t_0\ X for M\ F\ t_0 and X:: -
\Rightarrow - \Rightarrow - :: { linorder-topology}
sublocale martingale-order \subseteq submartingale using martingale-property by (unfold-locales)
(force simp add: integrable)+
```

 $sublocale martingale-linorder \subseteq submartingale-linorder ...$

7.4 Supermartingale

```
locale \ supermartingale = sigma-finite-adapted-process-order + sigma-finite-adapte
   assumes integrable: \bigwedge i. t_0 \leq i \Longrightarrow integrable \ M(Xi)
          and supermartingale-property: \bigwedge i \ j. \ t_0 \leq i \Longrightarrow i \leq j \Longrightarrow AE \ \xi \ in \ M. \ X \ i \ \xi
\geq cond\text{-}exp\ M\ (F\ i)\ (X\ j)\ \xi
locale supermartingale-linorder = supermartingale M F t_0 X for M F t_0 and X
:: - \Rightarrow - \Rightarrow - :: \{linorder-topology\}
sublocale martingale-order \subseteq supermartingale using martingale-property by (unfold-locales)
(force simp add: integrable)+
sublocale martingale-linorder \subseteq supermartingale-linorder ...
A stochastic process is a martingale, if and only if it is both a submartingale
and a supermartingale.
lemma martingale-iff:
   shows martingale M F t_0 X \longleftrightarrow submartingale M F t_0 X \land supermartingale M
F t_0 X
proof (rule iffI)
   assume asm: martingale M F t_0 X
   interpret martingale-order M F t_0 X by (intro martingale-order.intro asm)
    show submartingale M F t_0 X \wedge supermartingale M F t_0 X using submartin-
gale-axioms supermartingale-axioms by blast
   assume asm: submartingale M F t_0 X \wedge supermartingale M F t_0 X
   interpret submartingale\ M\ F\ t_0\ X by (simp\ add:\ asm)
   interpret supermartingale M F t_0 X by (simp add: asm)
  show martingale M F to X using submartingale-property supermartingale-property
by (unfold-locales) (intro integrable, blast, force)
\mathbf{qed}
7.5
              Martingale Lemmas
context martingale
begin
lemma cond-exp-diff-eq-zero:
   assumes t_0 \leq i \ i \leq j
   shows AE \xi in M. cond-exp M (F i) (\lambda \xi. X j \xi - X i \xi) \xi = 0
   using martingale-property[OF assms] assms
              sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i]
                  sigma-finite-subalgebra.cond-exp-diff[OF-integrable(1,1), of Fiji] by
fast force
lemma set-integral-eq:
   assumes A \in F \ i \ t_0 \le i \ i \le j
   shows set-lebesgue-integral M A (X i) = set-lebesgue-integral M A (X j)
proof -
   interpret sigma-finite-subalgebra M F i using assms(2) by blast
```

```
have \int x \in A. X \ i \ x \ \partial M = \int x \in A. cond-exp M (F \ i) (X \ j) x \ \partial M using
martingale-property[OF assms(2,3)] borel-measurable-cond-exp' assms subalgebras
subalgebra-def by (intro\ set\ -lebesgue\ -integral\ -cong\ -AE[OF\ -\ random\ -variable])\ fast-five formula of the subalgebra of the
force+
     also have ... = \int x \in A. X j x \partial M using assms by (auto simp: integrable intro:
cond\text{-}exp\text{-}set\text{-}integral[symmetric])
      finally show ?thesis.
qed
lemma scaleR-const[intro]:
     shows martingale M F t_0 (\lambda i \ x. \ c *_R X i \ x)
proof -
      {
            fix i j :: 'b assume asm: t_0 \le i \ i \le j
            interpret sigma-finite-subalgebra M F i using asm by blast
               have AE \ x \ in \ M. \ c *_R \ X \ i \ x = cond\text{-}exp \ M \ (F \ i) \ (\lambda x. \ c *_R \ X \ j \ x) \ x us-
ing asm cond-exp-scaleR-right[OF integrable, of j, THEN AE-symmetric] martin-
gale-property[OF asm] by force
     thus ?thesis by (unfold-locales) (auto simp add: integrable martingale.integrable)
qed
lemma uminus[intro]:
      shows martingale M F t_0 (-X)
      using scaleR-const[of -1] by (force\ intro:\ back-subst[of\ martingale\ M\ F\ t_0])
lemma add[intro]:
      assumes martingale M F t_0 Y
      shows martingale M F t_0 (\lambda i \xi. X i \xi + Y i \xi)
     interpret Y: martingale M F t_0 Y by (rule assms)
            fix i j :: 'b assume asm: t_0 \leq i i \leq j
            hence AE \xi in M. X i \xi + Y i \xi = cond\text{-}exp M (F i) (<math>\lambda x. X j x + Y j x) \xi
             \textbf{using } \textit{sigma-finite-subalgebra}. \textit{cond-exp-add}[\textit{OF-integrable } \textit{martingale.integrable}] \textit{OF-integrable } \textit{martingale.integrable}[\textit{OF-integrable } \textit{martingale.integrable}] \textit{OF-integrable } \textit{martingale.integrable}[\textit{OF-integrable } \textit{martingale.integrable}] \textit{OF-integrable}[\textit{OF-integrable } \textit{martingale.integrable}] \textit{OF-integrable}[\textit{OF-integrable}] \textit{Martingale.integrable}[\textit{OF-integrable}] \textit{Martingale.integrable}[\textit{OF-integrable}]
assms], of F i j j, THEN AE-symmetric]
                                        martingale	ext{-}property[OF\ asm]\ martingale	ext{-}martingale	ext{-}property[OF\ assms]
asm] by force
      }
      thus ?thesis using assms
     by (unfold-locales) (auto simp add: integrable martingale.integrable)
\mathbf{qed}
lemma diff[intro]:
      assumes martingale M F t_0 Y
      shows martingale M F t_0 (\lambda i x. X i x - Y i x)
     interpret Y: martingale M F t_0 Y by (rule assms)
```

```
fix i j :: 'b assume asm: t_0 \le i \ i \le j
   hence AE \xi in M. X i \xi - Y i \xi = cond\text{-}exp M (F i) (<math>\lambda x. X j x - Y j x) \xi
    \textbf{using } \textit{sigma-finite-subalgebra}. \textit{cond-exp-diff} [\textit{OF-integrable martingale.integrable} [\textit{OF-integrable martingale.integrable}] \\
assms], of F i j j, THEN AE-symmetric]
            martingale-property[OF asm] martingale.martingale-property[OF assms
asm] \ \mathbf{by} \ \mathit{fastforce}
  }
 thus ?thesis using assms by (unfold-locales) (auto simp add: integrable martin-
gale.integrable)
qed
end
lemma (in sigma-finite-adapted-process) martingale-of-cond-exp-diff-eq-zero:
  assumes integrable: \bigwedge i. t_0 \leq i \Longrightarrow integrable \ M \ (X \ i)
     and diff-zero: \bigwedge i \ j. t_0 \le i \Longrightarrow i \le j \Longrightarrow AE \ x \ in \ M. cond-exp M (F \ i) (\lambda \xi).
X j \xi - X i \xi) x = 0
   shows martingale M F t_0 X
proof
   fix i j :: 'b assume asm: t_0 \leq i i \leq j
   thus AE \xi in M. X i \xi = cond\text{-}exp M (F i) (X j) \xi
    using diff-zero [OF\ asm]\ sigma-finite-subalgebra.cond-exp-diff[OF\ -integrable(1,1),
of F i j i
          sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] by
fastforce
qed (intro integrable)
lemma (in sigma-finite-adapted-process) martingale-of-set-integral-eq:
  assumes integrable: \bigwedge i. t_0 \leq i \Longrightarrow integrable \ M(Xi)
     and \bigwedge A \ i \ j. \ t_0 \le i \Longrightarrow i \le j \Longrightarrow A \in F \ i \Longrightarrow set-lebesgue-integral \ M \ A \ (X)
i) = set-lebesgue-integral M A (X j)
   shows martingale M F t_0 X
proof (unfold-locales)
 fix i j :: 'b assume asm: t_0 \le i \ i \le j
 interpret sigma-finite-subalgebra M F i using asm by blast
 interpret r: sigma-finite-measure restr-to-subalg M (F i) by (simp add: sigma-fin-subalg)
  {
   fix A assume A \in restr-to-subalg M (F i)
   hence *: A \in F i using sets-restr-to-subalg subalgebras asm by blast
  have set-lebesgue-integral (restr-to-subalg M(Fi)) A(Xi) = set-lebesgue-integral
MA(Xi) using * subalq asm by (auto simp: set-lebesque-integral-def intro: inte-
gral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator)
    also have ... = set-lebesgue-integral M A (cond-exp M (F i) (X j)) using *
assms(2)[OF\ asm]\ cond-exp-set-integral[OF\ integrable]\ asm\ {f by}\ auto
  finally have set-lebesque-integral (restr-to-subalg M(Fi)) A(Xi) = set-lebesque-integral
(restr-to-subalq\ M\ (F\ i))\ A\ (cond-exp\ M\ (F\ i)\ (X\ j))\ using * subalq\ by\ (auto\ simp:
set-lebesque-integral-def intro!: integral-subalgebra2 [symmetric] borel-measurable-scaleR
```

```
borel-measurable-cond-exp borel-measurable-indicator)
       hence AE \xi in restr-to-subalg M (F i). X i \xi = cond\text{-}exp M (F i) (X j) \xi us-
ing asm by (intro r.density-unique-banach, auto intro: integrable-in-subalg subalg
borel-measurable-cond-exp integrable)
      thus AE \xi in M. Xi \xi = cond\text{-}exp\ M\ (Fi)\ (Xj)\ \xi using AE\text{-}restr\text{-}to\text{-}subalg[OF]
subalg] by blast
qed (simp add: integrable)
                               Submartingale Lemmas
7.6
{f context} submartingale
begin
lemma cond-exp-diff-nonneg:
       assumes t_0 \leq i \ i \leq j
       shows AE x in M. cond-exp M (F i) (\lambda \xi. X j \xi - X i \xi) x \ge 0
     using submartingale-property[OF assms] assms sigma-finite-subalgebra.cond-exp-diff[OF
-\ integrable(1,1),\ of\ -\ j\ i]\ sigma-finite-subalgebra.cond-exp-F-meas[OF\ -\ integrable
adapted, of i by fastforce
lemma add[intro]:
       assumes submartingale M F t_0 Y
       shows submartingale M F t_0 (\lambda i \xi. X i \xi + Y i \xi)
       interpret Y: submartingale M F t_0 Y by (rule assms)
        {
               fix i j :: 'b assume asm: t_0 \le i \ i \le j
               hence AE \xi in M. X i \xi + Y i \xi \leq cond\text{-}exp M (F i) (<math>\lambda x. X j x + Y j x) \xi
                {\bf using} \ sigma-finite-subalgebra. cond-exp-add [OF-integrable \ submarting ale. integrable [OF-integrable \ submarting ale. integrable \ submarting ale. integrable \ submarting ale. integrable \ submarting ale. integrable \ submarting \ submarting
assms], of F i j j]
                                      submartingale-property[OF asm] submartingale-submartingale-property[OF
assms asm] add-mono[of X i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i - Y i -
    thus ?thesis using assms by (unfold-locales) (auto simp add: borel-measurable-add
random-variable adapted integrable Y.random-variable Y.adapted submartingale.integrable)
ged
lemma diff[intro]:
       assumes supermartingale M F t_0 Y
       shows submartingale M F t_0 (\lambda i \xi. X i \xi - Y i \xi)
proof -
       interpret Y: supermartingale M F t_0 Y by (rule assms)
               fix i j :: 'b assume asm: t_0 \leq i \ i \leq j
               hence AE \xi in M. X i \xi - Y i \xi \leq cond\text{-}exp M (F i) (<math>\lambda x. X j x - Y j x) \xi
                \textbf{using } \textit{sigma-finite-subalgebra}. cond-exp-diff [\textit{OF-integrable } \textit{supermartingale.integrable}] \textit{OF-integrable } \textit{supermartingale.integrable} [\textit{OF-integrable } \textit{supermartingale.integrable}] \textit{OF-integrable } \textit{supermartingale.integrable} \textit{oF-integrable} \textit{oF-integr
assms], of F i j j]
```

```
submartingale-property[OF asm] supermartingale-supermartingale-property[OF
assms asm] diff-mono[ of X i - - - Y i - ] by force
 thus ?thesis using assms by (unfold-locales) (auto simp add: borel-measurable-diff
random-variable adapted integrable Y.random-variable Y.adapted supermartingale.integrable)
qed
lemma scaleR-nonneg:
 assumes c \geq \theta
 shows submartingale M F t_0 (\lambda i \xi. c *_R X i \xi)
proof
  {
   fix i j :: 'b assume asm: t_0 \le i \ i \le j
   thus AE \xi in M. c *_R X i \xi \leq cond\text{-}exp M (F i) (<math>\lambda \xi. c *_R X j \xi) \xi
      using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of F i
j\ c]\ submartingale\mbox{-property}[OF\ asm]\ {f by}\ (fastforce\ intro!:\ scaleR\mbox{-left-mono}[OF\ -
assms])
 }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)
lemma scaleR-le-zero:
 assumes c \leq \theta
 shows supermartingale M F t_0 (\lambda i \xi. c *_R X i \xi)
proof
   fix i j :: 'b assume asm: t_0 \le i \ i \le j
   thus AE \xi in M. c *_R X i \xi \ge cond\text{-}exp M (F i) (<math>\lambda \xi. c *_R X j \xi) \xi
     using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of F i j
c] submartingale-property[OF asm]
          by (fastforce intro!: scaleR-left-mono-neg[OF - assms])
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)
lemma uminus[intro]:
 shows supermartingale M F t_0 (-X)
  unfolding fun-Compl-def using scaleR-le-zero [of -1] by simp
end
context submartingale-linorder
begin
lemma set-integral-le:
 assumes A \in F \ i \ t_0 \le i \ i \le j
 shows set-lebesgue-integral M A (X i) \leq set-lebesgue-integral M A (X j)
 using submartingale-property[OF assms(2), of j] assms subalgebras
```

```
(auto\ intro!:\ scaleR-left-mono\ integral-mono-AE-banach\ integrable-mult-indicator
integrable simp add: subalgebra-def set-lebesgue-integral-def)
lemma max:
  assumes submartingale-linorder M F t_0 Y
  shows submartingale-linorder M F t_0 (\lambda i \ \xi. max (X i \ \xi) (Y i \ \xi))
proof (unfold-locales)
  interpret Y: submartingale-linorder M F t_0 Y by (rule assms)
   fix i j :: 'b assume asm: t_0 \le i i \le j
    have AE \ \xi \ in \ M. \ max \ (X \ i \ \xi) \ (Y \ i \ \xi) \le max \ (cond-exp \ M \ (F \ i) \ (X \ j) \ \xi)
(cond\text{-}exp\ M\ (F\ i)\ (Y\ j)\ \xi) using submartingale\text{-}property\ Y.submartingale\text{-}property
asm unfolding max-def by fastforce
   thus AE \notin in M. max(X i \notin)(Y i \notin) < cond-exp M(F i)(\lambda \notin. max(X j \notin)(Y i \notin))
(j \xi)) \xi using sigma-finite-subalgebra.cond-exp-max[OF - integrable Y.integrable, of
F \ i \ j \ j] \ asm \ \mathbf{by} \ (fast \ intro: \ order.trans)
  show \bigwedge i. t_0 \leq i \Longrightarrow (\lambda \xi. \ max \ (X \ i \ \xi) \ (Y \ i \ \xi)) \in borel-measurable \ (F \ i) \bigwedge i.
t_0 \leq i \implies integrable \ M \ (\lambda \xi. \ max \ (X \ i \ \xi) \ (Y \ i \ \xi)) \ \mathbf{by} \ (force \ intro: \ Y.integrable
integrable \ assms)+
qed
lemma max-\theta:
  shows submartingale-linorder M F t_0 (\lambda i \xi. max \theta (X i \xi))
  interpret zero: martingale-linorder M F t_0 \lambda- -. 0 by (force intro: martin-
gale-linorder.intro martingale-order.intro)
 show ?thesis by (intro zero.max submartingale-linorder.intro submartingale-axioms)
qed
end
lemma (in sigma-finite-adapted-process-order) submartingale-of-cond-exp-diff-nonneg:
  assumes integrable: \bigwedge i. t_0 \leq i \implies integrable M(X i)
      and diff-nonneg: \bigwedge i \ j. t_0 \le i \Longrightarrow i \le j \Longrightarrow AE \ x \ in \ M. cond-exp M (F \ i)
(\lambda \xi. \ X \ j \ \xi - X \ i \ \xi) \ x \ge 0
   shows submartingale M F t_0 X
proof (unfold-locales)
  {
   fix i j :: 'b assume asm: t_0 \le i \ i \le j
   thus AE \xi in M. X i \xi \leq cond\text{-}exp M (F i) (X j) \xi
       using diff-nonneg[OF asm] sigma-finite-subalgebra.cond-exp-diff[OF - inte-
grable(1,1), of F i j i
           sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] by
fast force
qed (intro integrable)
```

by (subst sigma-finite-subalgebra.cond-exp-set-integral OF - integrable assms(1),

of j])

```
\mathbf{lemma} \ (\mathbf{in} \ sigma-finite-adapted-process-linorder) \ submartingale-of-set-integral-le:
  assumes integrable: \bigwedge i. t_0 \leq i \Longrightarrow integrable \ M(Xi)
     and \bigwedge A \ i \ j. \ t_0 \leq i \Longrightarrow i \leq j \Longrightarrow A \in F \ i \Longrightarrow set-lebesgue-integral M \ A \ (X)
i) \leq set-lebesgue-integral M \land (X \ j)
   shows submartingale M F t_0 X
proof (unfold-locales)
   fix i j :: 'b assume asm: t_0 \le i \ i \le j
  interpret r: sigma-finite-measure restr-to-subalg M (Fi) using asm sigma-finite-subalgebra.sigma-fin-subalg
by blast
     fix A assume A \in restr-to-subalq M (F i)
     hence *: A \in F i using asm sets-restr-to-subalg subalgebras by blast
   have set-lebesque-integral (restr-to-subaly M(F_i)) A(X_i) = set-lebesque-integral
M A (X i) using * asm subalgebras by (auto simp: set-lebesque-integral-def intro:
integral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator)
      also have ... \leq set-lebesgue-integral M A (cond-exp M (F i) (X j)) using
* assms(2)[OF \ asm] asm \ sigma-finite-subalgebra.cond-exp-set-integral[OF - inte-
grable] by fastforce
     also have ... = set-lebesque-integral (restr-to-subalq M (F i)) A (cond-exp M
(F i) (X j) using * asm subalgebras by (auto simp: set-lebesgue-integral-def intro!:
integral-subalgebra2[symmetric] \ borel-measurable-scaleR \ borel-measurable-cond-exp
borel-measurable-indicator)
    finally have 0 \le set-lebesgue-integral (restr-to-subalg M (F i)) A (\lambda \xi. cond-exp
M(F i)(X j) \xi - X i \xi) using * asm subalgebras by (subst set-integral-diff,
auto simp add: set-integrable-def sets-restr-to-subalg intro!: integrable adapted inte-
grable-in-subalg\ borel-measurable-scale R\ borel-measurable-indicator\ borel-measurable-cond-exp
integrable-mult-indicator)
   }
   hence AE \xi in restr-to-subalg M (F i). 0 \leq cond-exp M (F i) (X j) \xi - X i \xi
   by (intro r.density-nonneq integrable-in-subalq asm subalqebras borel-measurable-diff
borel-measurable-cond-exp adapted Bochner-Integration integrable-diff integrable-cond-exp
integrable)
  thus AE \xi in M. Xi \xi \leq cond\text{-}exp M (Fi) (Xj) \xi using AE\text{-}restr\text{-}to\text{-}subalg[OF]
subalgebras asm by simp
  }
qed (intro integrable)
7.7
       Supermartingale Lemmas
The following lemmas are exact duals of the ones for submartingales.
context supermartingale
begin
lemma cond-exp-diff-nonneg:
 assumes t_0 \leq i \ i \leq j
 shows AE x in M. cond-exp M (F i) (\lambda \xi. X i \xi - X j \xi) x \ge 0
 \textbf{using} \ assms \ supermarting a le-property [OF \ assms] \ sigma-finite-subalgebra. cond-exp-diff [OF \ assms]
```

```
- integrable(1,1), of F i i j
                                       sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] by
fastforce
lemma add[intro]:
        assumes supermartingale M F t_0 Y
        shows supermartingale M F t_0 (\lambda i \xi. X i \xi + Y i \xi)
        interpret Y: supermartingale M F t_0 Y by (rule assms)
         {
                fix i j :: 'b assume asm: t_0 \leq i \ i \leq j
                hence AE \xi in M. X i \xi + Y i \xi \geq cond\text{-}exp M (F i) (<math>\lambda x. X j x + Y j x) \xi
                 \textbf{using } \textit{sigma-finite-subalgebra}. cond-exp-add [\textit{OF-integrable } \textit{supermartingale.integrable}] \textit{OF-integrable } \textit{supermartingale.integrable} [\textit{OF-integrable } \textit{supermartingale.integrable}] \textit{OF-integrable } \textit{supermartingale.integrable} \textit{oF-integrable } \textit{supermartingale.integrable} \textit{oF-integrable} \textit{oF-integrable
assms], of F i j j]
                                   supermarting a le-property [OF\ asm]\ supermarting a le-supermarting a le-property [OF\ asm]\ supermarting a le-property [OF\ asm]\ supermarting
assms asm] add-mono[of - X i - - Y i -] by force
     thus ?thesis using assms by (unfold-locales) (auto simp add: borel-measurable-add
random-variable adapted integrable Y-random-variable Y-adapted supermartingale.integrable)
qed
lemma diff[intro]:
        assumes submartingale M F t_0 Y
        shows supermartingale M F t_0 (\lambda i \xi. X i \xi - Y i \xi)
proof -
        interpret Y: submartingale M F t_0 Y by (rule assms)
                fix i j :: 'b assume asm: t_0 \leq i \ i \leq j
                hence AE \xi in M. X i \xi - Y i \xi \geq cond\text{-}exp M (F i) (<math>\lambda x. X j x - Y j x) \xi
                 \textbf{using } \textit{sigma-finite-subalgebra.} \textit{cond-exp-diff} [\textit{OF-integrable } \textit{submartingale.} \textit{integrable} [\textit{OF} \textit{-integrable } \textit{-integrable } \textit{submartingale.} \textit{-integrable } \textit
assms], of F i j j, unfolded fun-diff-def]
                                   supermarting a le-property [OF\ asm]\ submarting a le-submarting a le-property [OF\ asm]
assms asm] diff-mono[of - X i - Y i -] by force
     thus ?thesis using assms by (unfold-locales) (auto simp add: borel-measurable-diff
random-variable adapted integrable Y.random-variable Y.adapted submartingale.integrable)
qed
lemma scaleR-nonneg:
        assumes c \geq \theta
        shows supermartingale M F t_0 (\lambda i \xi. c *_R X i \xi)
proof
         {
                fix i j :: 'b assume asm: t_0 \le i \ i \le j
                thus AE \xi in M. c *_R X i \xi \geq cond\text{-}exp M (F i) (\lambda \xi. c *_R X j \xi) \xi
                               using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of F i
j c] supermartingale-property[OF asm] by (fastforce intro!: scaleR-left-mono[OF -
```

```
assms])
{\bf qed} \ (auto \ simp \ add: \ borel-measurable-integrable \ borel-measurable-scaleR \ integrable
random-variable adapted borel-measurable-const-scaleR)
\mathbf{lemma}\ scaleR-le-zero:
 assumes c \leq \theta
 shows submartingale M F t_0 (\lambda i \xi. c *_R X i \xi)
proof
   fix i j :: 'b assume asm: t_0 \leq i \ i \leq j
   thus AE \xi in M. c *_R X i \xi \leq cond\text{-}exp M (F i) (<math>\lambda \xi. c *_R X j \xi) \xi
    \mathbf{using} \ sigma-finite\text{-}subalgebra.cond\text{-}exp\text{-}scaleR\text{-}right[OF\text{-}integrable,\ of\ F\ i\ j\ c]
supermartingale-property[OF asm] by (fastforce intro!: scaleR-left-mono-neg[OF -
assms)
 }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)
lemma uminus[intro]:
 shows submartingale M F t_0 (-X)
 unfolding fun-Compl-def using scaleR-le-zero[of -1] by simp
end
context supermartingale-linorder
begin
lemma set-integral-ge:
 assumes A \in F \ i \ t_0 \le i \ i \le j
 shows set-lebesgue-integral M A (X i) \geq set-lebesgue-integral M A (X j)
 using supermartingale-property [OF\ assms(2),\ of\ j]\ assms\ subalgebras
  by (subst sigma-finite-subalgebra.cond-exp-set-integral OF - integrable assms(1),
   (auto\ intro!:\ scaleR-left-mono\ integral-mono-AE-banach\ integrable-mult-indicator
integrable simp add: subalgebra-def set-lebesque-integral-def)
lemma min:
 assumes supermartingale-linorder M F t_0 Y
 shows supermartingale-linorder M F t_0 (\lambda i \xi. min (X i \xi) (Y i \xi))
proof (unfold-locales)
 interpret Y: supermartingale-linorder M F t_0 Y by (rule assms)
   fix i j :: 'b assume asm: t_0 \leq i \ i \leq j
  have AE \xi in M. min(X i \xi)(Y i \xi) \ge min(cond-exp M(F i)(X j)\xi)(cond-exp)
M(Fi)(Yj)\xi) using supermartingale-property Y.supermartingale-property asm
unfolding min-def by fastforce
   thus AE \xi in M. min(X i \xi)(Y i \xi) \geq cond\text{-}exp M(F i)(\lambda \xi. min(X j \xi)(Y i \xi))
j \xi)) \xi using sigma-finite-subalgebra.cond-exp-min[OF - integrable Y.integrable, of
```

```
F \ i \ j \ j] \ asm \ \mathbf{by} \ (fast \ intro: \ order.trans)
  show \bigwedge i. t_0 \leq i \Longrightarrow (\lambda \xi. \ min \ (X \ i \ \xi) \ (Y \ i \ \xi)) \in borel-measurable \ (F \ i) \ \bigwedge i.
t_0 \leq i \implies integrable \ M \ (\lambda \xi. \ min \ (X \ i \ \xi) \ (Y \ i \ \xi)) \  by (force intro: Y.integrable
integrable \ assms)+
qed
lemma min-\theta:
  shows supermartingale-linorder M F t_0 (\lambda i \xi. min \theta (X i \xi))
  interpret zero: martingale-linorder M F t_0 \lambda- -. 0 by (force intro: martin-
gale-linorder.intro)
   show ?thesis by (intro zero.min supermartingale-linorder.intro supermartin-
gale-axioms)
qed
end
lemma (in sigma-finite-adapted-process-order) supermartingale-of-cond-exp-diff-le-zero:
 assumes integrable: \bigwedge i. t_0 \leq i \Longrightarrow integrable \ M \ (X \ i)
      and diff-le-zero: \bigwedge i \ j. \ t_0 \le i \Longrightarrow i \le j \Longrightarrow AE \ x \ in \ M. \ cond-exp \ M \ (F \ i)
(\lambda \xi. \ X \ j \ \xi - X \ i \ \xi) \ x \le 0
    shows supermartingale M F t_0 X
proof
    fix i j :: 'b assume asm: t_0 \le i \ i \le j
    thus AE \xi in M. X i \xi \geq cond\text{-}exp M (F i) (X j) \xi
       using diff-le-zero[OF asm] sigma-finite-subalgebra.cond-exp-diff[OF - inte-
grable(1,1), of F i j i
           sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] by
fast force
qed (intro integrable)
lemma (in sigma-finite-adapted-process-linorder) supermartingale-of-set-integral-ge:
  assumes integrable: \bigwedge i. t_0 \leq i \Longrightarrow integrable \ M \ (X \ i)
      and \bigwedge A \ i \ j. \ t_0 \leq i \Longrightarrow i \leq j \Longrightarrow A \in F \ i \Longrightarrow set\text{-lebesgue-integral} \ M \ A \ (X)
j) \leq set-lebesgue-integral M \land (X \mid i)
    shows supermartingale M F t_0 X
proof -
  interpret -: adapted-process M F t_0 - X by (rule uminus-adapted)
  interpret uminus-X: sigma-finite-adapted-process-linorder M F t_0 -X ..
 \mathbf{note} * = set\text{-}integral\text{-}uminus[unfolded\ set\text{-}integrable\text{-}def,\ OF\ integrable\text{-}mult\text{-}indicator[OF\ ]}
- integrable]]
  have supermartingale M F t_0 (-(-X))
   \mathbf{using} \ ord\text{-}eq\text{-}le\text{-}trans[OF* ord\text{-}le\text{-}eq\text{-}trans[OF le\text{-}imp\text{-}neg\text{-}le[OF assms(2)]*}[symmetric]]]
    by (intro submartingale.uminus uminus-X.submartingale-of-set-integral-le)
       (clarsimp simp add: fun-Compl-def subalgebra-def integrable | fastforce)+
```

```
thus ?thesis unfolding fun-Compl-def by simp qed
```

7.8 Discrete Time Martingales

```
locale nat-martingale = martingale M F 0 :: nat X for M F X
locale nat-submartingale = submartingale M F 0 :: nat X for M F X
locale nat-supermartingale = supermartingale M F 0 :: nat X for M F X
locale nat-submartingale-linorder = submartingale-linorder M F 0 :: nat X for M
FX
locale nat-supermartingale-linorder = supermartingale-linorder M F \theta :: nat X
for M F X
\mathbf{sublocale}\ nat\text{-}submartingale\text{-}linorder\subseteq nat\text{-}submartingale ..
sublocale nat-supermartingale-linorder \subseteq nat-supermartingale ...
lemma (in nat-martingale) predictable-const:
 assumes nat-predictable-process M F X
 shows AE \xi in M. X i \xi = X j \xi
proof -
 have *: AE \xi in M. X i \xi = X \theta \xi  for i
 proof (induction i)
   \mathbf{case}\ \theta
   then show ?case by (simp add: bot-nat-def)
  next
   case (Suc\ i)
  interpret S: nat-adapted-process M F \lambda i. X (Suc i) by (intro nat-predictable-process.adapted-Suc
   show ?case using Suc S.adapted[of i] martingale-property[OF - le-SucI, of i]
sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable, of F i Suc i] by fastforce
 show ?thesis using *[of i] *[of j] by force
qed
lemma (in nat-sigma-finite-adapted-process) martingale-of-set-integral-eq-Suc:
 assumes integrable: \bigwedge i. integrable M(X i)
    and \bigwedge A \ i.\ A \in F \ i \Longrightarrow set\text{-lebesgue-integral} \ M \ A \ (X \ i) = set\text{-lebesgue-integral}
M A (X (Suc i))
   shows nat-martingale M F X
proof (intro nat-martingale.intro martingale-of-set-integral-eq)
  fix i j A assume asm: i \leq j A \in sets (F i)
  show set-lebesgue-integral M A (X i) = set-lebesgue-integral M A (X j) using
asm
 proof (induction j - i arbitrary: i j)
   then show ?case using asm by simp
 next
   case (Suc\ n)
```

```
hence *: n = j - Suc \ i \ \mathbf{by} \ linarith
       have Suc\ i \leq j using Suc(2,3) by linarith
        thus ?case using sets-F-mono[OF - le\text{-}SucI] Suc(4) Suc(1)[OF *] by (auto
intro: assms(2)[THEN trans])
   ged
qed (simp add: integrable)
lemma (in nat-sigma-finite-adapted-process) martingale-nat:
    assumes integrable: \bigwedge i. integrable M(X i)
          and \bigwedge i. AE \xi in M. X i \xi = cond-exp M (F i) (X (Suc i)) \xi
       shows nat-martingale M F X
proof (unfold-locales)
   fix i j :: nat assume asm: i \leq j
   show AE \xi in M. X i \xi = cond\text{-}exp M (F i) (X j) \xi using asm
   proof (induction j - i arbitrary: i j)
       case \theta
       hence j = i by simp
    thus ?case using sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted,
 THEN AE-symmetric] by blast
    next
       case (Suc \ n)
       have j: j = Suc (n + i) using Suc by linarith
       have n: n = n + i - i using Suc by linarith
      have *: AE \xi in M. cond\text{-}exp M (F (n + i)) (X j) \xi = X (n + i) \xi  unfolding
j using assms(2)[THEN\ AE-symmetric] by blast
       have AE \xi in M. cond-exp M (F i) (X j) \xi = cond-exp M (F i) (cond-exp M)
(F(n+i))(Xj) \xi by (intro cond-exp-nested-subalg integrable subalg, simp add:
subalgebra-def sets-F-mono)
       hence AE \xi in M. cond-exp M (F i) (X j) \xi = cond-exp M (F i) <math>(X (n + i))
\xi using cond-exp-cong-AE[OF integrable-cond-exp integrable *] by force
       thus ?case using Suc(1)[OF n] by fastforce
qed (simp add: integrable)
\mathbf{lemma} \ (\mathbf{in} \ nat\text{-}sigma\text{-}finite\text{-}adapted\text{-}process) \ martingale\text{-}of\text{-}cond\text{-}exp\text{-}diff\text{-}Suc\text{-}eq\text{-}zero:
   assumes integrable: \bigwedge i. integrable M(X i)
          and \bigwedge i. AE \xi in M. cond-exp M (F i) (\lambda \xi. X (Suc i) \xi – X i \xi) \xi = 0
       shows nat-martingale M F X
proof (intro martingale-nat integrable)
  show AE \xi in M. Xi \xi = cond\text{-}exp M (Fi) (X (Suc i)) \xi using cond\text{-}exp\text{-}diff[OF]
integrable(1,1), of i Suc i i] cond-exp-F-meas[OF integrable adapted, of i] assms(2)[of integrable(1,1), of i Suc i i] cond-exp-F-meas[OF integr
i] by fastforce
qed
```

7.9 Discrete Time Submartingales

lemma (in nat-submartingale) predictable-mono: assumes nat-predictable-process $M F X i \leq j$

```
shows AE \xi in M. X i \xi \leq X j \xi
  using assms(2)
proof (induction j - i arbitrary: i j)
  case \theta
  then show ?case by simp
next
  case (Suc \ n)
 hence *: n = j - Suc \ i by linarith
 interpret S: nat-adapted-process M F \lambda i. X (Suc i) by (intro nat-predictable-process.adapted-Suc
 have Suc\ i \leq j using Suc(2,3) by linarith
  thus ?case using Suc(1)[OF *] S.adapted[of i] submartingale-property[OF -
le-SucI, of i] sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable, of F i Suc
i] by fastforce
qed
lemma (in nat-sigma-finite-adapted-process-linorder) submartingale-of-set-integral-le-Suc:
 assumes integrable: \bigwedge i. integrable M(X i)
    and \bigwedge A \ i.\ A \in F \ i \Longrightarrow set-lebesgue-integral M \ A \ (X \ i) \le set-lebesgue-integral
M A (X (Suc i))
   shows nat-submartingale M F X
proof (intro nat-submartingale.intro submartingale-of-set-integral-le)
  fix i j A assume asm: i \leq j A \in sets (F i)
  show set-lebesgue-integral M A (X \ i) \leq set-lebesgue-integral M A (X \ j) using
asm
  proof (induction j - i arbitrary: i j)
   then show ?case using asm by simp
  next
   case (Suc \ n)
   hence *: n = j - Suc \ i \ by \ linarith
   have Suc\ i \leq j using Suc(2,3) by linarith
    thus ?case using sets-F-mono[OF - le-SucI] Suc(4) Suc(1)[OF *] by (auto
intro: assms(2)[THEN \ order-trans])
 qed
ged (simp add: integrable)
lemma (in nat-sigma-finite-adapted-process-linorder) submartingale-nat:
  assumes integrable: \bigwedge i. integrable M(X i)
     and \bigwedge i. AE \xi in M. X i \xi \leq cond\text{-}exp\ M\ (F\ i)\ (X\ (Suc\ i))\ \xi
   shows nat-submartingale M F X
 using subalg integrable assms(2)
 \textbf{by } (intro\ submarting a le-of-set-integral-le-Suc\ ord-le-eq-trans[OF\ set-integral-mono-AE-banach]) \\
cond-exp-set-integral[symmetric]], <math>simp)
   (meson in-mono integrable-mult-indicator set-integrable-def subalgebra-def, me-
son integrable-cond-exp in-mono integrable-mult-indicator set-integrable-def subal-
gebra-def, fast+)
```

 ${\bf lemma\ (in\ }nat\text{-}sigma\text{-}finite\text{-}adapted\text{-}process\text{-}linorder)\ submartingale\text{-}of\text{-}cond\text{-}exp\text{-}diff\text{-}Suc\text{-}nonneg:}$

```
shows nat-submartingale M F X
proof (intro submartingale-nat integrable)
   \mathbf{fix} i
  show AE \xi in M. X i \xi \leq cond\text{-}exp M (F i) (X (Suc i)) \xi using cond\text{-}exp\text{-}diff[OF]
integrable (1,1), of i Suc i i] cond-exp-F-meas[OF integrable adapted, of i] assms(2)[of
i by fastforce
qed
lemma (in nat-submartingale-linorder) partial-sum-scaleR:
   assumes nat-adapted-process M F C \wedge i. AE \xi in M. 0 \leq C i \xi \wedge i. AE \xi in
M. Ci \xi \leq R
   shows nat-submartingale M F (\lambda n \xi. \sum i < n. C i \xi *_R (X (Suc i) \xi - X i \xi))
proof-
   interpret C: nat-adapted-process M F C by (rule assms)
   interpret C': nat-adapted-process M F \lambda i \xi. C (i-1) \xi *_R (X i \xi - X (i-1) \xi)
- 1) ξ) by (intro nat-adapted-process.intro adapted-process.scaleR-right-adapted
adapted-process.diff-adapted, unfold-locales) (auto intro: adaptedD C.adaptedD)+
   interpret C'': nat-adapted-process M F \lambda n \xi. \sum i < n. C i \xi *_R (X (Suc i) \xi - i)
X \ i \ \xi) by (rule C'.partial-sum-Suc-adapted[unfolded diff-Suc-1])
   interpret S: nat-sigma-finite-adapted-process-linorder M F (\lambda n \xi. \sum i < n. C i \xi
*_R (X (Suc i) \xi - X i \xi)) ..
  have integrable M (\lambda x. C i x *_R (X (Suc i) x - X i x)) for <math>i using assms(2,3)[of
i] by (intro Bochner-Integration.integrable-bound[OF integrable-scaleR-right, OF
Bochner-Integration.integrable-diff, OF integrable (1,1), of R Suc i i) (auto simp
add: mult-mono)
   moreover have AE \xi in M. 0 \leq cond\text{-}exp M (F i) (\lambda \xi. (\sum i < Suc i. C i \xi *_R
(X (Suc i) \xi - X i \xi)) - (\sum i < i. C i \xi *_R (X (Suc i) \xi - X i \xi))) \xi for i
         {\bf using} \ \ sigma-finite-subalgebra. cond-exp-measurable-scale R[OF-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calculation-calcul
 C.adapted, of i
                cond-exp-diff-nonneg[OF - le-SucI, OF - order.refl, of i] assms(2,3)[of\ i]
by (fastforce simp add: scaleR-nonneg-nonneg integrable)
   ultimately show ?thesis by (intro S.submartingale-of-cond-exp-diff-Suc-nonneg
Bochner-Integration.integrable-sum, blast+)
qed
lemma (in nat-submartingale-linorder) partial-sum-scaleR':
   assumes nat-predictable-process M F C \wedge i. AE \xi in M. 0 \leq C i \xi \wedge i. AE \xi in
M. Ci \xi \leq R
   shows nat-submartingale M F (\lambda n \xi. \sum i < n. C (Suc i) \xi *_R (X (Suc i) \xi - X)
i \xi)
proof
   {\bf interpret} \ C: \ nat\text{-}predictable\text{-}process \ M \ F \ C \ {\bf by} \ (rule \ assms)
  interpret Suc-C: nat-adapted-process M F \lambda i. C (Suc i) using C.adapted-Suc.
   show ?thesis by (intro partial-sum-scaleR[of - R] assms) (intro-locales)
qed
```

and $\bigwedge i$. AE ξ in M. cond-exp M (F i) ($\lambda \xi$. X (Suc i) $\xi - X$ i ξ) $\xi \geq 0$

assumes integrable: $\bigwedge i$. integrable M(X i)

7.10 Discrete Time Supermartingales

```
lemma (in nat-supermartingale) predictable-mono:
 assumes nat-predictable-process M F X i \leq j
 shows AE \xi in M. X i \xi \geq X j \xi
 using assms(2)
proof (induction j - i arbitrary: i j)
  case \theta
  then show ?case by simp
next
  case (Suc\ n)
 hence *: n = j - Suc \ i by linarith
 interpret S: nat-adapted-process M F \lambda i. X (Suc i) by (intro nat-predictable-process.adapted-Suc
 have Suc\ i \leq j using Suc(2,3) by linarith
  thus ?case using Suc(1)[OF *] S.adapted[of i] supermartingale-property[OF -]
le-SucI, of i] sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable, of F i Suc
i by fastforce
qed
\mathbf{lemma} \ (\mathbf{in} \ nat\text{-}sigma\text{-}finite\text{-}adapted\text{-}process\text{-}linorder) \ supermarting a le-of\text{-}set\text{-}integral\text{-}ge\text{-}Suc:
  assumes integrable: \bigwedge i. integrable M(X i)
    and \bigwedge A \ i.\ A \in F \ i \Longrightarrow set-lebesgue-integral \ M \ A \ (X \ i) \ge set-lebesgue-integral
M A (X (Suc i))
   shows nat-supermartingale M F X
proof -
 interpret -: adapted-process M F \theta - X by (rule uminus-adapted)
 interpret uminus-X: nat-sigma-finite-adapted-process-linorder M F - X...
 \mathbf{note} * = set\text{-}integral\text{-}uminus[unfolded set\text{-}integrable\text{-}def, OF integrable\text{-}mult\text{-}indicator]OF
- integrable]]
 have nat-supermartingale M F (-(-X))
  using ord-eq-le-trans[OF*ord-le-eq-trans[OF\ le-imp-neq-le[OF\ assms(2)]*[symmetric]]]
subalgebras
  by (intro nat-supermartingale.intro submartingale.uminus nat-submartingale.axioms
uminus-X.submartingale-of-set-integral-le-Suc)
      (clarsimp simp add: fun-Compl-def subalgebra-def integrable | fastforce)+
 thus ?thesis unfolding fun-Compl-def by simp
qed
\mathbf{lemma} \ (\mathbf{in} \ \mathit{nat-sigma-finite-adapted-process-linorder}) \ \mathit{supermartingale-nat}:
 assumes integrable: \bigwedge i. integrable M(X i)
     and \bigwedge i. AE \xi in M. X i \xi \geq cond\text{-}exp\ M\ (F\ i)\ (X\ (Suc\ i))\ \xi
   shows nat-supermartingale M F X
proof -
 interpret -: adapted-process M F \theta - X by (rule uminus-adapted)
 interpret uminus-X: nat-sigma-finite-adapted-process-linorder MF-X...
 have AE \xi in M. - X i \xi \leq cond\text{-}exp M (F i) (\lambda x. - X (Suc i) x) \xi for i using
assms(2) cond-exp-uminus[OF integrable, of i Suc i] by force
  hence nat-supermartingale M F (-(-X)) by (intro nat-supermartingale intro
submarting a le. uminus\ nat-submarting a le. axioms\ uminus-X. submarting a le-nat)\ (auto
```

```
simp add: fun-Compl-def integrable)
 thus ?thesis unfolding fun-Compl-def by simp
qed
lemma (in nat-sigma-finite-adapted-process-linorder) supermartingale-of-cond-exp-diff-Suc-le-zero:
 assumes integrable: \bigwedge i. integrable M(X i)
     and \bigwedge i. AE \xi in M. cond-exp M (F i) (\lambda \xi. X (Suc i) \xi - X i \xi) \xi \leq 0
   shows nat-supermartingale M F X
proof (intro supermartingale-nat integrable)
 show AE \xi in M. Xi \xi \geq cond-exp M (Fi) (X (Suc i)) \xi using cond-exp-diff[OF]
integrable(1,1), of i Suc i i cond-exp-F-meas[OF integrable adapted, of i assms(2)] of
i] by fastforce
qed
end
theory Example-Coin-Toss
 imports Martingale HOL-Probability. Stream-Space HOL-Probability. Probability-Mass-Function
begin
8
     Example: Coin Toss
We consider a probability space consisting of infinite sequences of coin tosses.
definition bernoulli-stream :: real \Rightarrow (bool \ stream) measure where
  bernoulli-stream p = stream-space (measure-pmf (bernoulli-pmf p))
lemma space-bernoulli-stream[simp]: space (bernoulli-stream p) = UNIV by (simp
add: bernoulli-stream-def space-stream-space)
We define the fortune of the player at time n to be the number of heads
minus number of tails.
\mathbf{definition} \ \mathit{fortune} :: \mathit{nat} \Rightarrow \mathit{bool} \ \mathit{stream} \Rightarrow \mathit{real} \ \mathbf{where}
 fortune n = (\lambda s. \sum b \leftarrow stake (Suc n) s. if b then 1 else -1)
definition toss :: nat \Rightarrow bool stream \Rightarrow real where
  toss n = (\lambda s. if snth s n then 1 else -1)
lemma toss-indicator-def: toss n = indicator \{s. s !! n\} - indicator \{s. \neg s !! n\}
  unfolding toss-def indicator-def by force
lemma range-toss: range (toss \ n) = \{-1, 1\}
proof -
 have sconst True !! n by simp
 moreover have \neg sconst\ False\ !!\ n\ by\ simp
 ultimately have \exists x. x !! n \exists x. \neg x !! n  by blast+
 thus ?thesis unfolding toss-def image-def by auto
```

```
qed
```

```
lemma vimage-toss: toss n - A = (if 1 \in A \text{ then } \{s. s !! n\} \text{ else } \{\}) \cup (if - 1 \in A \text{ then } \{s. s !! n\} \text{ else } \{\})
A then \{s. \neg s !! n\} else \{\}\}
 unfolding vimage-def toss-def by auto
lemma fortune-Suc: fortune (Suc n) s = fortune \ n \ s + toss (Suc n) \ s
 by (induction n arbitrary: s) (simp add: fortune-def toss-def)+
lemma fortune-toss-sum: fortune n \ s = (\sum i \in \{..n\}. \ toss \ i \ s)
  by (induction n arbitrary: s) (simp add: fortune-def toss-def, simp add: for-
tune-Suc)
lemma fortune-bound: norm (fortune n \ s) \leq Suc \ n \ by (induction \ n) (force simp
add: fortune-toss-sum toss-def)+
Our definition of bernoulli-stream constitutes a probability space.
interpretation prob-space bernoulli-stream p unfolding bernoulli-stream-def by
(simp add: measure-pmf.prob-space-axioms prob-space.prob-space-stream-space)
abbreviation toss-filtration p \equiv nat-natural-filtration (bernoulli-stream p) toss
The stochastic process toss is adapted to the filtration it generates.
interpretation toss: nat-adapted-process bernoulli-stream p nat-natural-filtration
(bernoulli-stream p) toss toss
 by (intro nat-adapted-process.intro stochastic-process.adapted-process-natural-filtration)
    (unfold-locales, auto simp add: toss-def bernoulli-stream-def)
Similarly, the stochastic process fortune is adapted to the filtration generated
by the tosses.
interpretation fortune: nat-finite-adapted-process-linorder bernoulli-stream p nat-natural-filtration
(bernoulli-stream p) toss fortune
proof -
  show nat-finite-adapted-process-linorder (bernoulli-stream p) (toss-filtration p)
fortune
  unfolding fortune-toss-sum
  by (intro nat-finite-adapted-process-linorder.intro
          finite-adapted-process-linorder.intro
          finite-adapted-process-order.intro
          finite-adapted-process.intro
          toss.partial-sum-adapted[folded atMost-atLeast0]) intro-locales
qed
lemma integrable-toss: integrable (bernoulli-stream p) (toss n)
 using toss.random-variable
 by (intro Bochner-Integration.integrable-bound[OF integrable-const[of - 1 :: real]])
(auto simp add: toss-def)
```

lemma integrable-fortune: integrable (bernoulli-stream p) (fortune n) using fortune-bound

by $(intro\ Bochner-Integration.integrable-bound[OF\ integrable-const[of\ -\ Suc\ n]$ fortune.random-variable]) auto

We provide the following lemma to explicitly calculate the probability of events in this probability space.

```
\mathbf{lemma}\ \textit{measure-bernoulli-stream-snth-pred}\colon
    assumes 0 \le p and p \le 1 and finite J
    shows prob p \{ w \in space (bernoulli-stream p). \forall j \in J. P j = w !! j \} = p \cap card
(J \cap Collect \ P) * (1 - p) \cap (card \ (J - Collect \ P))
proof -
    let ?PiE = (\Pi_E \ i \in J. \ if \ P \ i \ then \ \{True\} \ else \ \{False\})
   have product-prob-space (\lambda-. measure-pmf (bernoulli-pmf p)) by unfold-locales
    hence *: to-stream - '\{s. \ \forall i \in J. \ P \ i = s \ !! \ i\} = \{s. \ \forall i \in J. \ P \ i = s \ i\} using
assms by (simp add: to-stream-def)
    also have ... = prod-emb UNIV (\lambda-. measure-pmf (bernoulli-pmf p)) J ?PiE
   proof -
       {
          fix s assume (\forall i \in J. P i = s i)
       hence (\forall i \in J. \ P \ i = s \ i) = (s \in prod\text{-}emb \ UNIV \ (\lambda\text{-}. measure\text{-}pmf \ (bernoulli\text{-}pmf \ )))
p)) \ J \ ?PiE)
                  by (subst prod-emb-iff[of s]) (smt (verit, best) not-def assms(3) id-def
PiE-eq-singleton UNIV-I extensional-UNIV insert-iff singletonD space-measure-pmf)
      moreover
       {
          fix s assume \neg(\forall i \in J. P i = s i)
          then obtain i where i \in J P i \neq s i by blast
       hence (\forall i \in J. \ P \ i = s \ i) = (s \in prod-emb\ UNIV\ (\lambda -. measure-pmf\ (bernoulli-pmf\ (
p)) \ J \ ?PiE)
           by (simp add: restrict-def prod-emb-iff[of s]) (smt (verit, ccfv-SIG) PiE-mem
assms(3) id-def insert-iff singleton-iff)
      ultimately show ?thesis by auto
   qed
    finally have inteq: (to-stream - '\{s. \forall i \in J. P \ i = s !! i\}) = prod-emb UNIV
(\lambda-. measure-pmf (bernoulli-pmf p)) J ?PiE.
   let ?M = (Pi_M \ UNIV \ (\lambda -. \ measure-pmf \ (bernoulli-pmf \ p)))
    have emeasure (bernoulli-stream p) \{s \in space (bernoulli-stream p), \forall i \in J. P i\}
= s \parallel i = emeasure ?M (to-stream - `\{s. \forall i \in J. P \mid i = s \parallel i \})
        using assms emeasure-distrof to-stream ?M (vimage-algebra (streams (space
(measure-pmf\ (bernoulli-pmf\ p))))\ (!!)\ ?M)\ \{s.\ \forall\ i\in J.\ P\ i=s\ !!\ i\},\ symmetric]
measurable-to-stream[of (measure-pmf (bernoulli-pmf p))]
      by (simp only: bernoulli-stream-def stream-space-def *, simp add: space-PiM)
(smt (verit, best) emeasure-notin-sets in-vimage-algebra inf-top.right-neutral sets-distr
vimage-Collect)
   also have ... = emeasure ?M (prod-emb UNIV (\lambda-. measure-pmf (bernoulli-pmf
```

```
p)) J ?PiE) using integ by (simp add: space-PiM)
 also have ... = (\prod i \in J. \text{ emeasure (measure-pmf (bernoulli-pmf p)) (if } P \text{ i then})
\{True\}\ else\ \{False\}))
  by (subst emeasure-PiM-emb) (auto simp add: prob-space-measure-pmf assms(3))
 also have ... = (\prod i \in J \cap Collect \ P. \ ennreal \ p) * (\prod i \in J - Collect \ P. \ ennreal \ p)
(1 - p)
  unfolding emeasure-pmf-single of bernoulli-pmf p True, unfolded pmf-bernoulli-True OF
assms(1,2)], symmetric]
        emeasure-pmf-single[of bernoulli-pmf p False, unfolded pmf-bernoulli-False[OF
assms(1,2)], symmetric]
   by (simp\ add:\ prod.Int-Diff[OF\ assms(3),\ of\ -\ Collect\ P])
 also have ... = p \land card (J \cap Collect P) * (1 - p) \land card (J - Collect P) using
assms by (simp add: prod-ennreal ennreal-mult' ennreal-power)
  finally show ?thesis using assms by (intro measure-eq-emeasure-eq-ennreal)
auto
qed
lemma
 assumes 0 \le p and p \le 1
 shows measure-bernoulli-stream-snth: prob p \{ w \in space (bernoulli-stream p). w \}
!! \ i\} = p
  and measure-bernoulli-stream-neg-snth: prob p \{ w \in space (bernoulli-stream p) \}.
\neg w !! i \} = 1 - p
 using measure-bernoulli-stream-snth-pred[OF assms, of \{i\} \lambda x. True]
      measure-bernoulli-stream-snth-pred [OF assms, of \{i\} \lambda x. False] by auto
Now we can express the expected value of a single coin toss.
lemma integral-toss:
 assumes 0 \le p \ p \le 1
 shows expectation p(toss n) = 2 * p - 1
proof -
 have [simp]: \{s. s !! n\} \in events p  using measurable-snth[THEN measurable-sets,
of {True} measure-pmf (bernoulli-pmf p) n, folded bernoulli-stream-def]
   by (simp add: vimage-def)
 have expectation \ p \ (toss \ n) = Bochner-Integration.simple-bochner-integral \ (bernoulli-stream
p) (toss n)
   using toss.random-variable[of n, THEN measurable-sets]
  by (intro simple-bochner-integrable-eq-integral [symmetric] simple-bochner-integrable.intros)
(auto simp add: toss-def simple-function-def image-def)
 also have ... = p - prob p \{s. \neg s !! n\} unfolding simple-bochner-integral-def
using measure-bernoulli-stream-snth[OF assms]
   by (simp add: range-toss, simp add: toss-def)
 also have ... = p - (1 - prob \ p \ \{s. \ s !! \ n\}) by (subst prob-compl[symmetric],
auto simp add: Collect-neg-eq Compl-eq-Diff-UNIV)
 finally show ?thesis using measure-bernoulli-stream-snth[OF assms] by simp
qed
```

Now, we show that the tosses are independent from one another.

 $\mathbf{lemma}\ indep ext{-}vars ext{-}toss:$

```
assumes 0 \le p \ p \le 1
  shows indep-vars p (\lambda-. borel) toss {0..}
proof (subst indep-vars-def, intro conjI indep-sets-sigma)
     fix A J assume asm: J \neq \{\} finite J \forall j \in J. A j \in \{toss \ j - `A \cap space\}
(bernoulli\text{-stream }p) \mid A. A \in borel\}
    hence \forall j \in J. \exists B \in borel. A j = toss j - `B \cap space (bernoulli-stream p) by
    then obtain B where B-is: A j = toss j - B j \cap space (bernoulli-stream p)
B j \in borel \ \mathbf{if} \ j \in J \ \mathbf{for} \ j \ \mathbf{by} \ met is
    have prob p \cap (A \cap J) = (\prod j \in J. \ prob \ p \cap (A \mid j))
    proof cases
We consider the case where there is a zero probability event.
      assume \exists j \in J. 1 \notin B j \land -1 \notin B j
      then obtain j where j-is: j \in J 1 \notin B j -1 \notin B j by blast
    hence A-j-empty: A j = \{\} using B-is by (force simp\ add: toss-def\ vimage-def)
      hence \bigcap (A \cdot J) = \{\} using j-is by blast
      moreover have prob p(A j) = 0 using A-j-empty by simp
      ultimately show ?thesis using j-is asm(2) by auto
    \mathbf{next}
We now assume all events have positive probability.
      assume \neg(\exists j \in J. \ 1 \notin B \ j \land -1 \notin B \ j)
      hence *: 1 \in B \ j \lor -1 \in B \ j \ \text{if} \ j \in J \ \text{for} \ j \ \text{using} \ that \ \text{by} \ blast
      define J' where [simp]: J' = \{j \in J. (1 \in B \ j) \longleftrightarrow (-1 \notin B \ j)\}
      hence toss \ j \ w \in B \ j \longleftrightarrow (1 \in B \ j) = w \ !! \ j \ \textbf{if} \ j \in J' \ \textbf{for} \ w \ j \ \textbf{using} \ that
unfolding toss-def by simp
      hence (\bigcap (A 'J')) = \{w \in space (bernoulli-stream p). \forall j \in J'. (1 \in B j) = a \}
w \parallel j using B-is by force
      hence prob-J': prob p \ (\bigcap (A \ 'J')) = p \ \widehat{} \ card \ (J' \cap \{j. \ 1 \in B \ j\}) * (1 - j)
p) \cap card (J' - \{j. \ 1 \in B \ j\})
            using measure-bernoulli-stream-snth-pred[OF assms finite-subset[OF -
asm(2)], of J' \lambda j. 1 \in B j] by auto
The index set J' consists of the indices of all non-trivial events.
     have A-j-True: A j = \{w \in space (bernoulli-stream p), w !! j\} if j \in J' \cap \{j, j\}
1 \in B j for j
        using that by (auto simp add: toss-def B-is(1) split: if-splits)
      have A-j-False: A j = \{w \in space (bernoulli-stream p). \neg w !! j\} if j \in J'
\{j. \ 1 \in B \ j\}  for j
```

have A-j-top: A j = space (bernoulli-stream p) if $j \in J - J'$ for j using that

using that B-is **by** (auto simp add: toss-def)

* by (auto simp add: B-is toss-def)

```
hence \bigcap (A ' J) = \bigcap (A ' J') by auto
     hence prob p (\bigcap (A 'J)) = prob p (\bigcap (A 'J')) by presburger
     also have ... = (\prod j \in J' \cap \{j, 1 \in B \}). prob p(A j) * (\prod j \in J' - \{j, 1 \in B \})
B j \}. prob p (A j))
      by (simp only: prob-J' A-j-True A-j-False measure-bernoulli-stream-snth[OF]
assms measure-bernoulli-stream-neg-snth[OF assms] cong: prod.cong) simp
   also have ... = (\prod j \in J'. prob p(A j)) using asm(2) by (intro prod.Int-Diff[symmetric])
auto
     also have ... = (\prod j \in J'. prob p(A j)) * (\prod j \in J - J'. prob p(A j)) using
A-j-top prob-space by simp
      also have ... = (\prod j \in J. \ prob \ p \ (A \ j)) using asm(2) by (metis \ (no-types,
lifting) J'-def mem-Collect-eq mult.commute prod.subset-diff subsetI)
     finally show ?thesis.
   qed
  thus indep-sets p (\lambda i. {toss i - 'A \cap space (bernoulli-stream p) |A. A \in sets
borel\}) \{0..\} using measurable-sets[OF toss.random-variable]
 by (intro indep-setsI subsetI) fastforce
qed (simp, intro Int-stableI, simp, metis sets.Int vimage-Int)
The fortune of a player is a martingale (resp. sub- or supermartingale) with
respect to the filtration generated by the coin tosses.
theorem fortune-martingale:
 assumes p = 1/2
 shows nat-martingale (bernoulli-stream p) (toss-filtration p) fortune
 \mathbf{using}\ cond-exp-indep[OF\ for tune. subalg\ indep-set-natural-filtration\ integrable-toss,
OF\ zero\text{-}order(1)\ lessI\ indep-vars-toss,\ of\ p
       integral-toss assms
     \mathbf{by} \ (\mathit{intro}\ \mathit{fortune}. \mathit{martingale-of-cond-exp-diff-Suc-eq-zero}\ \mathit{integrable-fortune})
(force simp add: fortune-toss-sum)
theorem fortune-submartingale:
 assumes 1/2 \le p \ p \le 1
 shows nat-submartingale (bernoulli-stream p) (toss-filtration p) fortune
proof (intro fortune.submartingale-of-cond-exp-diff-Suc-nonneg integrable-fortune)
 \mathbf{fix} \ n
 show AE \xi in bernoulli-stream p. 0 \le cond-exp (bernoulli-stream p) (toss-filtration
(p \ n) \ (\lambda \xi. \ fortune \ (Suc \ n) \ \xi - fortune \ n \ \xi) \ \xi
  \mathbf{using}\ cond-exp-indep[OF\ for tune. subalg\ indep-set-natural-filtration\ integrable-toss,
OF\ zero-order(1)\ lessI\ indep-vars-toss,\ of\ p\ n
         integral-toss[of p Suc n] assms
   by (force simp add: fortune-toss-sum)
qed
theorem fortune-supermartingale:
 assumes 0 \le p \ p \le 1/2
  shows nat-supermartingale (bernoulli-stream p) (toss-filtration p) fortune
proof (intro fortune.supermartingale-of-cond-exp-diff-Suc-le-zero integrable-fortune)
 \mathbf{fix} \ n
```

```
show AE \ \xi in bernoulli-stream p.\ 0 \ge cond-exp (bernoulli-stream p) (toss-filtration p n) (\lambda \xi. fortune (Suc n) \xi — fortune n \xi) \xi using cond-exp-indep[OF fortune.subalg indep-set-natural-filtration integrable-toss, OF zero-order(1) less I indep-vars-toss, of p n] integral-toss[of p Suc n] assms by (force simp add: fortune-toss-sum) qed end
```