

On the Formalization of Martingales

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September 13, 2023

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```

theory Measure-Space-Supplement
  imports HOL-Analysis.Measure-Space
begin

```

1 Sigma Algebra Generated by a Family of Functions

```

definition family-vimage-algebra :: 'a set  $\Rightarrow$  ('a  $\Rightarrow$  'b) set  $\Rightarrow$  'b measure  $\Rightarrow$  'a
measure where
  family-vimage-algebra  $\Omega$   $S$   $M \equiv$  sigma  $\Omega$  ( $\bigcup f \in S. \{f - 'A \cap \Omega \mid A. A \in M\}$ )

```

```

lemma family-vimage-algebra-singleton: family-vimage-algebra  $\Omega$  { $f$ }  $M =$  vimage-algebra  $\Omega$   $f$   $M$  unfolding family-vimage-algebra-def vimage-algebra-def by simp

```

```

lemma
  shows sets-family-vimage-algebra: sets (family-vimage-algebra  $\Omega$   $S$   $M$ ) = sigma-sets
 $\Omega$  ( $\bigcup f \in S. \{f - 'A \cap \Omega \mid A. A \in M\}$ )
  and space-family-vimage-algebra[simp]: space (family-vimage-algebra  $\Omega$   $S$   $M$ ) =
 $\Omega$ 
  by (auto simp add: family-vimage-algebra-def sets-measure-of-conv space-measure-of-conv)

```

```

lemma measurable-family-vimage-algebra:
  assumes  $f \in S$   $f \in \Omega \rightarrow$  space  $M$ 
  shows  $f \in$  family-vimage-algebra  $\Omega$   $S$   $M \rightarrow_M M$ 
  using assms by (intro measurableI, auto simp add: sets-family-vimage-algebra)

```

```

lemma measurable-family-vimage-algebra-singleton:
  assumes  $f \in \Omega \rightarrow$  space  $M$ 
  shows  $f \in$  family-vimage-algebra  $\Omega$  { $f$ }  $M \rightarrow_M M$ 
  using assms measurable-family-vimage-algebra by blast

```

```

lemma measurable-family-iff-sets:
  shows ( $S \subseteq N \rightarrow_M M$ )  $\longleftrightarrow$   $S \subseteq$  space  $N \rightarrow$  space  $M \wedge$  family-vimage-algebra
(space  $N$ )  $S$   $M \subseteq N$ 
proof (standard, goal-cases)
  case 1
    hence subset:  $S \subseteq$  space  $N \rightarrow$  space  $M$  using measurable-space by fast
    have  $\{f - 'A \cap$  space  $N \mid A. A \in M\} \subseteq N$  if  $f \in S$  for  $f$  using measurable-iff-sets[unfolded family-vimage-algebra-singleton[symmetric], of  $f$ ] 1 subset that
by (fastforce simp add: sets-family-vimage-algebra)
    then show ?case unfolding sets-family-vimage-algebra using sets.sigma-algebra-axioms
by (simp add: subset, intro sigma-algebra.sigma-sets-subset, blast+)
next
  case 2
    hence subset:  $S \subseteq$  space  $N \rightarrow$  space  $M$  by simp
    show ?case

```

```

proof (standard, goal-cases)
  case (1 x)
    have family-vimage-algebra (space N) {x} M  $\subseteq$  N by (metis (no-types, lifting)
1 2 sets-family-vimage-algebra SUP-le-iff sigma-sets-le-sets-iff singletonD)
    thus ?case using measurable-iff-sets[unfolded family-vimage-algebra-singleton[symmetric]]
subset[THEN subsetD, OF 1] by fast
  qed
qed

end
theory Elementary-Metric-Spaces-Supplement
  imports HOL-Analysis.Elementary-Metric-Spaces
begin

```

2 Diameter Lemma

```

lemma diameter-comp-strict-mono:
  fixes s :: nat  $\Rightarrow$  'a :: metric-space
  assumes strict-mono r bounded {s i | i. r n  $\leq$  i}
  shows diameter {s (r i) | i. n  $\leq$  i}  $\leq$  diameter {s i | i. r n  $\leq$  i}
proof (rule diameter-subset)
  show {s (r i) | i. n  $\leq$  i}  $\subseteq$  {s i | i. r n  $\leq$  i} using assms(1) monotoneD
strict-mono-mono by fastforce
qed (intro assms(2))

lemma diameter-bounded-bound':
  fixes S :: 'a :: metric-space set
  assumes S: bdd-above (case-prod dist ' (S $\times$ S)) x  $\in$  S y  $\in$  S
  shows dist x y  $\leq$  diameter S
proof -
  have (x,y)  $\in$  S $\times$ S using S by auto
  then have dist x y  $\leq$  (SUP (x,y) $\in$ S $\times$ S. dist x y) by (rule cSUP-upper2[OF
assms(1)]) simp
  with  $\langle x \in S \rangle$  show ?thesis by (auto simp: diameter-def)
qed

lemma bounded-imp-dist-bounded:
  assumes bounded (range s)
  shows bounded (( $\lambda(i, j).$  dist (s i) (s j)) ' ({n..}  $\times$  {n..}))
  using bounded-dist-comp[OF bounded-fst bounded-snd, OF bounded-Times(1,1)[OF
assms(1,1)]] by (rule bounded-subset, force)

lemma cauchy-iff-diameter-tends-to-zero-and-bounded:
  fixes s :: nat  $\Rightarrow$  'a :: metric-space
  shows Cauchy s  $\longleftrightarrow$  (( $\lambda n.$  diameter {s i | i. i  $\geq$  n})  $\longrightarrow$  0  $\wedge$  bounded (range
s))
proof -
  have {s i | i. N  $\leq$  i}  $\neq$  {} for N by blast
  hence diameter-SUP: diameter {s i | i. N  $\leq$  i} = (SUP (i, j)  $\in$  {N..}  $\times$  {N..}).

```

```

dist (s i) (s j)) for N unfolding diameter-def by (auto intro!: arg-cong[of - - Sup])
show ?thesis
proof ((intro iffI) ; clarsimp)
  assume asm: Cauchy s
  have  $\exists N. \forall n \geq N. \text{norm } (\text{diameter } \{s\ i \mid i. n \leq i\}) < e$  if e-pos:  $e > 0$  for e
  proof -
    obtain N where dist-less:  $\text{dist } (s\ n) (s\ m) < (1/2) * e$  if  $n \geq N\ m \geq N$ 
    for n m using asm e-pos by (metis Cauchy-def mult-pos-pos zero-less-divide-iff
    zero-less-numeral zero-less-one)
    {
      fix r assume  $r \geq N$ 
      hence  $\text{dist } (s\ n) (s\ m) < (1/2) * e$  if  $n \geq r\ m \geq r$  for n m using dist-less
    that by simp
      hence  $(\text{SUP } (i, j) \in \{r.. \} \times \{r.. \}. \text{dist } (s\ i) (s\ j)) \leq (1/2) * e$  by (intro
    cSup-least) fastforce+
      also have ... < e using e-pos by simp
      finally have  $\text{diameter } \{s\ i \mid i. r \leq i\} < e$  using diameter-SUP by presburger
    }
    moreover have  $\text{diameter } \{s\ i \mid i. r \leq i\} \geq 0$  for r unfolding diameter-SUP
    using bounded-imp-dist-bounded[OF cauchy-imp-bounded, THEN bounded-imp-bdd-above,
    OF asm] by (intro cSup-upper2, auto)
    ultimately show ?thesis by auto
  qed
  thus  $(\lambda n. \text{diameter } \{s\ i \mid i. n \leq i\}) \longrightarrow 0 \wedge \text{bounded } (\text{range } s)$  using
  cauchy-imp-bounded[OF asm] by (simp add: LIMSEQ-iff)
next
  assume asm:  $(\lambda n. \text{diameter } \{s\ i \mid i. n \leq i\}) \longrightarrow 0$  bounded (range s)
  have  $\exists N. \forall n \geq N. \forall m \geq N. \text{dist } (s\ n) (s\ m) < e$  if e-pos:  $e > 0$  for e
  proof -
    obtain N where diam-less:  $\text{diameter } \{s\ i \mid i. r \leq i\} < e$  if  $r \geq N$  for r
    using LIMSEQ-D asm(1) e-pos by fastforce
    {
      fix n m assume  $n \geq N\ m \geq N$ 
      hence  $\text{dist } (s\ n) (s\ m) < e$  using cSUP-lessD[OF bounded-imp-dist-bounded[THEN
    bounded-imp-bdd-above], OF asm(2) diam-less[unfolded diameter-SUP]] by auto
    }
    thus ?thesis by blast
  qed
  then show Cauchy s by (simp add: Cauchy-def)
qed
qed
end

```

```

theory Bochner-Integration-Supplement
imports HOL-Analysis.Bochner-Integration HOL-Analysis.Set-Integral Elementary-Metric-Spaces-Supplement

```

begin

3 Integrable Simple Functions

lemma *integrable-implies-simple-function-sequence:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{banach, second-countable-topology}\}$

assumes *integrable* $M f$

obtains s **where** $\bigwedge i. \text{simple-function } M (s\ i)$

and $\bigwedge i. \text{emeasure } M \{y \in \text{space } M. s\ i\ y \neq 0\} \neq \infty$

and $\bigwedge x. x \in \text{space } M \implies (\lambda i. s\ i\ x) \longrightarrow f\ x$

and $\bigwedge x\ i. x \in \text{space } M \implies \text{norm } (s\ i\ x) \leq 2 * \text{norm } (f\ x)$

proof –

have $f: f \in \text{borel-measurable } M (\int^+ x. \text{norm } (f\ x) \partial M) < \infty$ **using** *assms*
unfolding *integrable-iff-bounded* **by** *auto*

obtain s **where** $s: \bigwedge i. \text{simple-function } M (s\ i) \bigwedge x. x \in \text{space } M \implies (\lambda i. s\ i\ x) \longrightarrow f\ x \bigwedge i x. x \in \text{space } M \implies \text{norm } (s\ i\ x) \leq 2 * \text{norm } (f\ x)$ **using**
borel-measurable-implies-sequence-metric[*OF* $f(1)$] **unfolding** *norm-conv-dist* **by**
metis

{

fix i

have $(\int^+ x. \text{norm } (s\ i\ x) \partial M) \leq (\int^+ x. \text{ennreal } (2 * \text{norm } (f\ x)) \partial M)$ **using**
s **by** (*intro nn-integral-mono, auto*)

also have $\dots < \infty$ **using** f **by** (*simp add: nn-integral-cmult ennreal-mult-less-top ennreal-mult*)

finally have *sbi: Bochner-Integration.simple-bochner-integrable* $M (s\ i)$ **using**
s **by** (*intro simple-bochner-integrableI-bounded*) *auto*

hence $\text{emeasure } M \{y \in \text{space } M. s\ i\ y \neq 0\} \neq \infty$ **by** (*auto intro: integrableI-simple-bochner-integrable simple-bochner-integrable.cases*)

}

thus *?thesis* **using** *that s* **by** *blast*

qed

lemma *simple-function-indicator-representation:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$

assumes $f: \text{simple-function } M f$ **and** $x: x \in \text{space } M$

shows $f\ x = (\sum y \in f\ ' \text{space } M. \text{indicator } (f\ -\ ' \{y\} \cap \text{space } M) x *_R y)$

(**is** $?l = ?r$)

proof –

have $?r = (\sum y \in f\ ' \text{space } M.$

$(\text{if } y = f\ x \text{ then indicator } (f\ -\ ' \{y\} \cap \text{space } M) x *_R y \text{ else } 0))$ **by** (*auto intro!: sum.cong*)

also have $\dots = \text{indicator } (f\ -\ ' \{f\ x\} \cap \text{space } M) x *_R f\ x$ **using** *assms* **by** (*auto dest: simple-functionD*)

also have $\dots = f\ x$ **using** x **by** (*auto simp: indicator-def*)

finally show *?thesis* **by** *auto*

qed

lemma *simple-function-indicator-representation-AE:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$

assumes f : *simple-function* M f
shows $\text{AE } x \text{ in } M. f\ x = (\sum y \in f\text{' } \text{space } M. \text{indicator } (f - \text{' } \{y\} \cap \text{space } M) \ x$
 $\ast_R\ y)$
by (*metis* (*mono-tags*, *lifting*) *AE-I2 simple-function-indicator-representation* f)

lemmas *simple-function-scaleR*[*intro*] = *simple-function-compose2*[**where** $h = (\ast_R)$]

lemmas *integrable-simple-function* = *simple-bochner-integrable.intros*[*THEN has-bochner-integral-simple-bochner*
THEN integrable.intros]

lemma *integrable-simple-function-induct*[*consumes 2*, *case-names cong indicator*
add, *induct set: simple-function*]:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes f : *simple-function* M f *emeasure* M $\{y \in \text{space } M. f\ y \neq 0\} \neq \infty$
assumes *cong*: $\bigwedge f\ g. \text{simple-function } M\ f \implies \text{emeasure } M\ \{y \in \text{space } M. f\ y \neq 0\} \neq \infty$
 $\implies \text{simple-function } M\ g \implies \text{emeasure } M\ \{y \in \text{space } M. g\ y \neq 0\} \neq \infty$
 $\implies (\bigwedge x. x \in \text{space } M \implies f\ x = g\ x) \implies P\ f \implies P\ g$
assumes *indicator*: $\bigwedge A\ y. A \in \text{sets } M \implies \text{emeasure } M\ A < \infty \implies P\ (\lambda x. \text{indicator } A\ x \ast_R\ y)$
assumes *add*: $\bigwedge f\ g. \text{simple-function } M\ f \implies \text{emeasure } M\ \{y \in \text{space } M. f\ y \neq 0\} \neq \infty \implies$
 $\text{simple-function } M\ g \implies \text{emeasure } M\ \{y \in \text{space } M. g\ y \neq 0\} \neq \infty \implies$
 $(\bigwedge z. z \in \text{space } M \implies \text{norm } (f\ z + g\ z) = \text{norm } (f\ z) + \text{norm } (g\ z)) \implies$
 $P\ f \implies P\ g \implies P\ (\lambda x. f\ x + g\ x)$

shows $P\ f$

proof–

let $?f = \lambda x. (\sum y \in f\text{' } \text{space } M. \text{indicat-real } (f - \text{' } \{y\} \cap \text{space } M) \ x \ast_R\ y)$
have $f\text{-ae-eq}$: $f\ x = ?f\ x$ **if** $x \in \text{space } M$ **for** x **using** *simple-function-indicator-representation*[*OF*
 $f(1)$ *that*].

moreover have *emeasure* M $\{y \in \text{space } M. ?f\ y \neq 0\} \neq \infty$ **by** (*metis* (*no-types*,
lifting) *Collect-cong calculation* $f(2)$)

moreover have $P\ (\lambda x. \sum y \in S. \text{indicat-real } (f - \text{' } \{y\} \cap \text{space } M) \ x \ast_R\ y)$
 $\text{simple-function } M\ (\lambda x. \sum y \in S. \text{indicat-real } (f - \text{' } \{y\} \cap \text{space } M) \ x$
 $\ast_R\ y)$

emeasure M $\{y \in \text{space } M. (\sum x \in S. \text{indicat-real } (f - \text{' } \{x\} \cap \text{space } M) \ y \ast_R\ x) \neq 0\} \neq \infty$

if $S \subseteq f\text{' } \text{space } M$ **for** S **using** *simple-functionD*(1)[*OF assms*(1),
THEN rev-finite-subset, OF that] *that*

proof (*induction rule: finite-induct*)

case *empty*

{

case 1

then show $?case$ **using** *indicator*[*of* $\{\}$] **by force**

next

case 2
 then show ?case by force
 next
 case 3
 then show ?case by force
 }
 next
 case (insert x F)
 have $(f - \{x\} \cap \text{space } M) \subseteq \{y \in \text{space } M. f y \neq 0\}$ if $x \neq 0$ using that by blast
 moreover have $\{y \in \text{space } M. f y \neq 0\} = \text{space } M - (f - \{0\} \cap \text{space } M)$ by blast
 moreover have $\text{space } M - (f - \{0\} \cap \text{space } M) \in \text{sets } M$ using simple-functionD(2)[OF f(1)] by blast
 ultimately have fin-0: $\text{emeasure } M (f - \{x\} \cap \text{space } M) < \infty$ if $x \neq 0$ using that by (metis emeasure-mono f(2) infinity-enreal-def top.not-eq-extremum top-unique)
 hence fin-1: $\text{emeasure } M \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\} \neq \infty$ if $x \neq 0$ by (metis (mono-tags, lifting) emeasure-mono f(1) indicator-simps(2) linorder-not-less mem-Collect-eq scaleR-eq-0-iff simple-functionD(2) subsetI that)

 have *: $(\sum y \in \text{insert } x F. \text{indicat-real } (f - \{y\} \cap \text{space } M) xa *_R y) = (\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) xa *_R y) + \text{indicat-real } (f - \{x\} \cap \text{space } M) xa *_R x$ for xa by (metis (no-types, lifting) Diff-empty Diff-insert0 add.commute insert.hyps(1) insert.hyps(2) sum.insert-remove)
 have **: $\{y \in \text{space } M. (\sum x \in \text{insert } x F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \subseteq \{y \in \text{space } M. (\sum x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\}$ unfolding * by fastforce
 {
 case 1
 hence $x: x \in f - \text{space } M$ and $F: F \subseteq f - \text{space } M$ by auto
 show ?case
 proof (cases $x = 0$)
 case True
 then show ?thesis unfolding * using insert(3)[OF F] by simp
 next
 case False
 have norm-argument: $\text{norm } ((\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) + \text{indicat-real } (f - \{x\} \cap \text{space } M) z *_R x) = \text{norm } (\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) + \text{norm } (\text{indicat-real } (f - \{x\} \cap \text{space } M) z *_R x)$
 if $z: z \in \text{space } M$ for z
 proof (cases $f z = x$)
 case True
 have $\text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y = 0$ if $y \in F$ for y using True insert(2) z that 1 unfolding indicator-def by force
 hence $(\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) = 0$ by (meson sum.neutral)
 }


```

    then show ?thesis by force
  next
    case False
    then show ?thesis by force
  qed
  show ?thesis using False simple-functionD(2)[OF f(1)] insert(3,5)[OF F]
  simple-function-scaleR fin-0 fin-1 by (subst *, subst add, subst simple-function-sum)
  (blast intro: norm-argument indicator)+
  qed
next
  case 2
  hence x:  $x \in f \text{ ` space } M$  and F:  $F \subseteq f \text{ ` space } M$  by auto
  show ?case
  proof (cases  $x = 0$ )
    case True
    then show ?thesis unfolding * using insert(4)[OF F] by simp
  next
    case False
    then show ?thesis unfolding * using insert(4)[OF F] simple-functionD(2)[OF
f(1)] by fast
  qed
next
  case 3
  hence x:  $x \in f \text{ ` space } M$  and F:  $F \subseteq f \text{ ` space } M$  by auto
  show ?case
  proof (cases  $x = 0$ )
    case True
    then show ?thesis unfolding * using insert(5)[OF F] by simp
  next
    case False
    have emeasure M  $\{y \in \text{space } M. (\sum_{x \in \text{insert } x F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \leq \text{emeasure } M \{y \in \text{space } M. (\sum_{x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\}$ 
    using ** simple-functionD(2)[OF insert(4)[OF F]] simple-functionD(2)[OF
f(1)] by (intro emeasure-mono, force+)
    also have ...  $\leq \text{emeasure } M \{y \in \text{space } M. (\sum_{x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} + \text{emeasure } M \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\}$ 
    using simple-functionD(2)[OF insert(4)[OF F]] simple-functionD(2)[OF
f(1)] by (intro emeasure-subadditive, force+)
    also have ...  $< \infty$  using insert(5)[OF F] fin-1[OF False] by (simp add:
less-top)
    finally show ?thesis by simp
  qed
}
qed
moreover have simple-function M  $(\lambda x. \sum_{y \in f \text{ ` space } M. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y)$  using calculation by blast

```

moreover have $P (\lambda x. \sum_{y \in f^{-1} \text{ space } M} \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_{\mathbb{R}} y)$ **using** *calculation by blast*
ultimately show *?thesis* **by** (*intro cong[OF - - f(1,2)]*, *blast*, *presburger+*)
qed

lemma *integrable-simple-function-induct-nn*[*consumes 3*, *case-names cong indicator add*, *induct set: simple-function*]:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes f : *simple-function* M f *emeasure* $M \{y \in \text{space } M. f y \neq 0\} \neq \infty \wedge x. x \in \text{space } M \longrightarrow f x \geq 0$

assumes *cong*: $\wedge f g. \text{simple-function } M f \implies \text{emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty \implies (\wedge x. x \in \text{space } M \implies f x \geq 0) \implies \text{simple-function } M g \implies \text{emeasure } M \{y \in \text{space } M. g y \neq 0\} \neq \infty \implies (\wedge x. x \in \text{space } M \implies g x \geq 0) \implies (\wedge x. x \in \text{space } M \implies f x = g x) \implies P f \implies P g$

assumes *indicator*: $\wedge A y. y \geq 0 \implies A \in \text{sets } M \implies \text{emeasure } M A < \infty \implies P (\lambda x. \text{indicator } A x *_{\mathbb{R}} y)$

assumes *add*: $\wedge f g. (\wedge x. x \in \text{space } M \implies f x \geq 0) \implies \text{simple-function } M f \implies \text{emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty \implies$

$(\wedge x. x \in \text{space } M \implies g x \geq 0) \implies \text{simple-function } M g \implies \text{emeasure } M \{y \in \text{space } M. g y \neq 0\} \neq \infty \implies$

$(\wedge z. z \in \text{space } M \implies \text{norm } (f z + g z) = \text{norm } (f z) + \text{norm } (g z)) \implies$

$P f \implies P g \implies P (\lambda x. f x + g x)$

shows $P f$

proof–

let $?f = \lambda x. (\sum_{y \in f^{-1} \text{ space } M} \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_{\mathbb{R}} y)$

have $f\text{-ae-eq}$: $f x = ?f x$ **if** $x \in \text{space } M$ **for** x **using** *simple-function-indicator-representation[OF f(1) that]* .

moreover have $\text{emeasure } M \{y \in \text{space } M. ?f y \neq 0\} \neq \infty$ **by** (*metis (no-types, lifting) Collect-cong calculation f(2)*)

moreover have $P (\lambda x. \sum_{y \in S} \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_{\mathbb{R}} y)$
 $\text{simple-function } M (\lambda x. \sum_{y \in S} \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_{\mathbb{R}} y)$

$\text{emeasure } M \{y \in \text{space } M. (\sum_{x \in S} \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_{\mathbb{R}} x) \neq 0\} \neq \infty$

$\wedge x. x \in \text{space } M \implies 0 \leq (\sum_{y \in S} \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_{\mathbb{R}} y)$

if $S \subseteq f^{-1} \text{ space } M$ **for** S **using** *simple-functionD(1)[OF assms(1), THEN rev-finite-subset, OF that]* **that**

proof (*induction rule: finite-subset-induct'*)

case *empty*

{

case *1*

then show *?case* **using** *indicator[of 0 {}]* **by force**

next

case *2*

then show *?case* **by force**

```

next
  case 3
  then show ?case by force
next
  case 4
  then show ?case by force
}
next
  case (insert x F)
  have  $(f - \{x\} \cap \text{space } M) \subseteq \{y \in \text{space } M. f y \neq 0\}$  if  $x \neq 0$  using that by
blast
  moreover have  $\{y \in \text{space } M. f y \neq 0\} = \text{space } M - (f - \{0\} \cap \text{space } M)$ 
by blast
  moreover have  $\text{space } M - (f - \{0\} \cap \text{space } M) \in \text{sets } M$  using simple-functionD(2)[OF f(1)] by blast
  ultimately have fin-0:  $\text{emeasure } M (f - \{x\} \cap \text{space } M) < \infty$  if  $x \neq 0$ 
using that by (metis emeasure-mono f(2) infinity-ennreal-def top.not-eq-extremum
top-unique)
  hence fin-1:  $\text{emeasure } M \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\} \neq \infty$  if  $x \neq 0$  by (metis (mono-tags, lifting) emeasure-mono f(1) indicator-simps(2) linorder-not-less mem-Collect-eq scaleR-eq-0-iff simple-functionD(2) subsetI that)

  have nonneg-x:  $x \geq 0$  using insert f by blast
  have *:  $(\sum_{y \in \text{insert } x F. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y) = (\sum_{y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y) + \text{indicat-real } (f - \{x\} \cap \text{space } M) x *_R x$  for  $x$  by (metis (no-types, lifting) add.commute insert.hyps(1) insert.hyps(4) sum.insert)
  have **:  $\{y \in \text{space } M. (\sum_{x \in \text{insert } x F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \subseteq \{y \in \text{space } M. (\sum_{x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\}$  unfolding
* by fastforce
  {
    case 1
    show ?case
    proof (cases  $x = 0$ )
      case True
      then show ?thesis unfolding * using insert by simp
    next
      case False
      have norm-argument:  $\text{norm } ((\sum_{y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) + \text{indicat-real } (f - \{x\} \cap \text{space } M) z *_R x) = \text{norm } (\sum_{y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) + \text{norm } (\text{indicat-real } (f - \{x\} \cap \text{space } M) z *_R x)$ 
if  $z: z \in \text{space } M$  for  $z$ 
      proof (cases  $f z = x$ )
        case True
        have  $\text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y = 0$  if  $y \in F$  for  $y$  using
True insert z that 1 unfolding indicator-def by force
        hence  $(\sum_{y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) = 0$  by (meson

```

```

sum.neutral)
  thus ?thesis by force
  qed (force)
  show ?thesis using False fin-0 fin-1 f norm-argument by (subst *, subst add,
presburger add: insert, intro insert, intro insert, fastforce simp add: indicator-def
intro!: insert(2) f(3), auto intro!: indicator insert nonneg-x)
  qed
next
case 2
show ?case
proof (cases x = 0)
  case True
  then show ?thesis unfolding * using insert by simp
next
case False
  then show ?thesis unfolding * using insert simple-functionD(2)[OF f(1)]
by fast
qed
next
case 3
show ?case
proof (cases x = 0)
  case True
  then show ?thesis unfolding * using insert by simp
next
case False
  have emeasure M {y ∈ space M. (∑ x∈insert x F. indicat-real (f - ' {x}
  ∩ space M) y *R x) ≠ 0} ≤ emeasure M ({y ∈ space M. (∑ x∈F. indicat-real (f
  - ' {x} ∩ space M) y *R x) ≠ 0} ∪ {y ∈ space M. indicat-real (f - ' {x} ∩ space
  M) y *R x ≠ 0})
  using ** simple-functionD(2)[OF insert(6)] simple-functionD(2)[OF f(1)]
  insert.IH(2) by (intro emeasure-mono, blast, simp)
  also have ... ≤ emeasure M {y ∈ space M. (∑ x∈F. indicat-real (f - ' {x}
  ∩ space M) y *R x) ≠ 0} + emeasure M {y ∈ space M. indicat-real (f - ' {x} ∩
  space M) y *R x ≠ 0}
  using simple-functionD(2)[OF insert(6)] simple-functionD(2)[OF f(1)]
by (intro emeasure-subadditive, force+)
  also have ... < ∞ using insert(7) fin-1[OF False] by (simp add: less-top)
  finally show ?thesis by simp
qed
next
case (4 ξ)
  thus ?case using insert nonneg-x f(3) by (auto simp add: scaleR-nonneg-nonneg
intro: sum-nonneg)
}
qed
moreover have simple-function M (λx. ∑ y∈f ' space M. indicat-real (f - ' {y}
  ∩ space M) x *R y) using calculation by blast
moreover have P (λx. ∑ y∈f ' space M. indicat-real (f - ' {y} ∩ space M) x

```

```

*_R y) using calculation by blast
  moreover have  $\bigwedge x. x \in \text{space } M \implies 0 \leq f x$  using  $f(3)$  by simp
  ultimately show ?thesis by (intro cong[OF - - -  $f(1,2)$ ], blast, blast, fast)
presburger+
qed

lemma finite-nn-integral-imp-ae-finite:
  fixes  $f :: 'a \Rightarrow \text{ennreal}$ 
  assumes  $f \in \text{borel-measurable } M$   $(\int^+ x. f x \, \partial M) < \infty$ 
  shows  $AE \, x \text{ in } M. f x < \infty$ 
proof (rule ccontr, goal-cases)
  case 1
  let ?A =  $\text{space } M \cap \{x. f x = \infty\}$ 
  have *:  $\text{emeasure } M \, ?A > 0$  using 1 assms(1) by (metis (mono-tags, lifting)
    assms(2) eventually-mono infinity-ennreal-def nn-integral-not-eq-infinite top.not-eq-extremum)
  have  $(\int^+ x. f x * \text{indicator } ?A \, x \, \partial M) = (\int^+ x. \infty * \text{indicator } ?A \, x \, \partial M)$  by
    (metis (mono-tags, lifting) indicator-inter-arith indicator-simps(2) mem-Collect-eq
    mult-eq-0-iff)
  also have  $\dots = \infty * \text{emeasure } M \, ?A$  using assms(1) by (intro nn-integral-cmult-indicator,
    simp)
  also have  $\dots = \infty$  using * by fastforce
  finally have  $(\int^+ x. f x * \text{indicator } ?A \, x \, \partial M) = \infty$  .
  moreover have  $(\int^+ x. f x * \text{indicator } ?A \, x \, \partial M) \leq (\int^+ x. f x \, \partial M)$  by (intro
    nn-integral-mono, simp add: indicator-def)
  ultimately show ?case using assms(2) by simp
qed

```

```

lemma cauchy-L1-AE-cauchy-subseq:
  fixes  $s :: \text{nat} \Rightarrow 'a \Rightarrow 'b :: \{\text{banach, second-countable-topology}\}$ 
  assumes [measurable]:  $\bigwedge n. \text{integrable } M \, (s \, n)$ 
  and  $\bigwedge e. e > 0 \implies \exists N. \forall i \geq N. \forall j \geq N. \text{LINT } x | M. \text{norm } (s \, i \, x - s \, j \, x) < e$ 
  obtains  $r$  where strict-mono  $r \, AE \, x \text{ in } M. \text{Cauchy } (\lambda i. s \, (r \, i) \, x)$ 
proof-
  have  $\exists r. \forall n. (\forall i \geq r \, n. \forall j \geq r \, n. \text{LINT } x | M. \text{norm } (s \, i \, x - s \, j \, x) < (1 / 2) ^ n) \wedge (r \, (\text{Suc } n) > r \, n)$ 
  proof (intro dependent-nat-choice, goal-cases)
    case 1
    then show ?case using assms(2) by presburger
  next
    case (2  $x \, n$ )
    obtain  $N$  where *:  $\text{LINT } x | M. \text{norm } (s \, i \, x - s \, j \, x) < (1 / 2) ^ \text{Suc } n$  if  $i \geq N$ 
     $j \geq N$  for  $i \, j$  using assms(2)[of  $(1 / 2) ^ \text{Suc } n$ ] by auto
    {
      fix  $i \, j$  assume  $i \geq \max N (\text{Suc } x) \, j \geq \max N (\text{Suc } x)$ 
      hence  $\text{integral}^L M \, (\lambda x. \text{norm } (s \, i \, x - s \, j \, x)) < (1 / 2) ^ \text{Suc } n$  using * by
      fastforce
    }
  qed

```

$\}$
then show *?case by fastforce*
qed
then obtain r **where** *strict-mono: strict-mono r and $\forall i \geq r \ n. \forall j \geq r \ n. \text{LINT } x | M. \text{norm } (s \ i \ x - s \ j \ x) < (1 / 2) \wedge n$ for n using strict-mono-Suc-iff by blast*
hence r -is: *LINT $x | M. \text{norm } (s \ (r \ (\text{Suc } n)) \ x - s \ (r \ n) \ x) < (1 / 2) \wedge n$ for n by (simp add: strict-mono-leD)*

define g **where** $g = (\lambda n \ x. (\sum i \leq n. \text{ennreal } (\text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x))))$
define g' **where** $g' = (\lambda x. \sum i. \text{ennreal } (\text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x)))$

have *integrable-g: $(\int^+ x. g \ n \ x \ \partial M) < 2$ for n*
proof –
have $(\int^+ x. g \ n \ x \ \partial M) = (\int^+ x. (\sum i \leq n. \text{ennreal } (\text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x))) \ \partial M)$ **using** *g-def by simp*
also have $\dots = (\sum i \leq n. (\int^+ x. \text{ennreal } (\text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x)) \ \partial M))$ **by** *(intro nn-integral-sum, simp)*
also have $\dots = (\sum i \leq n. \text{LINT } x | M. \text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x))$
unfolding *dist-norm using assms(1) by (subst nn-integral-eq-integral[OF integrable-norm], auto)*
also have $\dots < \text{ennreal } (\sum i \leq n. (1 / 2) \wedge i)$ **by** *(intro ennreal-lessI[OF sum-pos sum-strict-mono[OF finite-atMost - r-is]], auto)*
also have $\dots \leq \text{ennreal } 2$ **unfolding** *sum-gp0[of 1 / 2 n] by (intro ennreal-leI, auto)*
finally show $(\int^+ x. g \ n \ x \ \partial M) < 2$ **by** *simp*
qed

have *integrable-g': $(\int^+ x. g' \ x \ \partial M) \leq 2$*
proof –
have *incseq $(\lambda n. g \ n \ x)$ for x by (intro incseq-SucI, auto simp add: g-def ennreal-leI)*
hence *convergent $(\lambda n. g \ n \ x)$ for x unfolding convergent-def using LIMSEQ-incseq-SUP by fast*
hence $(\lambda n. g \ n \ x) \longrightarrow g' \ x$ **for** x **unfolding** *g-def g'-def by (intro summable-iff-convergent'[THEN iffD2, THEN summable-LIMSEQ], blast)*
hence $(\int^+ x. g' \ x \ \partial M) = (\int^+ x. \liminf (\lambda n. g \ n \ x) \ \partial M)$ **by** *(metis lim-imp-Liminf trivial-limit-sequentially)*
also have $\dots \leq \liminf (\lambda n. \int^+ x. g \ n \ x \ \partial M)$ **by** *(intro nn-integral-liminf, simp add: g-def)*
also have $\dots \leq \liminf (\lambda n. 2)$ **using** *integrable-g by (intro Liminf-mono) (simp add: order-le-less)*
also have $\dots = 2$ **using** *sequentially-bot tendsto-iff-Liminf-eq-Limsup by blast*
finally show *?thesis .*
qed
hence *AE x in $M. g' \ x < \infty$ by (intro finite-nn-integral-imp-ae-finite) (auto simp add: order-le-less-trans g'-def)*
moreover have *summable $(\lambda i. \text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x))$ if $g' \ x \neq \infty$ for x using that unfolding g'-def by (intro summable-suminf-not-top) fastforce+*

ultimately have *ae-summable*: $AE\ x\ in\ M.\ summable\ (\lambda i.\ s\ (r\ (Suc\ i))\ x - s\ (r\ i)\ x)$ using *summable-norm-cancel* **unfolding** *dist-norm* **by** *force*

```
{
  fix x assume summable (λi. s (r (Suc i)) x - s (r i) x)
  hence (λn. ∑ i<n. s (r (Suc i)) x - s (r i) x) ⟶ (∑ i. s (r (Suc i)) x -
s (r i) x) using summable-LIMSEQ by blast
  moreover have (λn. (∑ i<n. s (r (Suc i)) x - s (r i) x)) = (λn. s (r n) x -
s (r 0) x) using sum-lessThan-telescope by fastforce
  ultimately have (λn. s (r n) x - s (r 0) x) ⟶ (∑ i. s (r (Suc i)) x - s
(r i) x) by argo
  hence (λn. s (r n) x - s (r 0) x + s (r 0) x) ⟶ (∑ i. s (r (Suc i)) x - s
(r i) x) + s (r 0) x by (intro isCont-tendsto-compose[of - λz. z + s (r 0) x], auto)
  hence Cauchy (λn. s (r n) x) by (simp add: LIMSEQ-imp-Cauchy)
}
hence AE x in M. Cauchy (λi. s (r i) x) using ae-summable by fast
thus ?thesis by (rule that[OF strict-mono(1)])
qed
```

4 Totally Ordered Banach Spaces

lemma *integrable-max*[*simp*, *intro*]:

```
fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology}
assumes fg[measurable]: integrable M f integrable M g
shows integrable M (λx. max (f x) (g x))
proof (rule Bochner-Integration.integrable-bound)
  {
    fix x y :: 'b
    have norm (max x y) ≤ max (norm x) (norm y) by linarith
    also have ... ≤ norm x + norm y by simp
    finally have norm (max x y) ≤ norm (norm x + norm y) by simp
  }
  thus AE x in M. norm (max (f x) (g x)) ≤ norm (norm (f x) + norm (g x)) by
simp
qed (auto intro: Bochner-Integration.integrable-add[OF integrable-norm[OF fg(1)]
integrable-norm[OF fg(2)])])
```

lemma *integrable-min*[*simp*, *intro*]:

```
fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology}
assumes [measurable]: integrable M f integrable M g
shows integrable M (λx. min (f x) (g x))
proof -
  have norm (min (f x) (g x)) ≤ norm (f x) ∨ norm (min (f x) (g x)) ≤ norm (g
x) for x by linarith
  thus ?thesis by (intro integrable-bound[OF integrable-max[OF integrable-norm(1,1),
OF assms]], auto)
qed
```

lemma *integral-nonneg-AE-banach*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes $[measurable]: f \in \text{borel-measurable } M$ **and** $\text{nonneg}: AE\ x\ \text{in } M. 0 \leq f\ x$

shows $0 \leq \text{integral}^L\ M\ f$

proof *cases*

assume *integrable*: $\text{integrable } M\ f$

hence $\text{max}: (\lambda x. \text{max } 0\ (f\ x)) \in \text{borel-measurable } M \wedge x. 0 \leq \text{max } 0\ (f\ x)$

integrable $M\ (\lambda x. \text{max } 0\ (f\ x))$ **by** *auto*

hence $0 \leq \text{integral}^L\ M\ (\lambda x. \text{max } 0\ (f\ x))$

proof $-$

obtain s **where** $*$: $\wedge i. \text{simple-function } M\ (s\ i)$

$\wedge i. \text{emeasure } M\ \{y \in \text{space } M. s\ i\ y \neq 0\} \neq \infty$

$\wedge x. x \in \text{space } M \implies (\lambda i. s\ i\ x) \longrightarrow f\ x$

$\wedge x\ i. x \in \text{space } M \implies \text{norm } (s\ i\ x) \leq 2 * \text{norm } (f\ x)$ **using**

integrable-implies-simple-function-sequence[*OF integrable*] **by** *blast*

have *simple*: $\wedge i. \text{simple-function } M\ (\lambda x. \text{max } 0\ (s\ i\ x))$ **using** $*$ **by** *fast*

have $\wedge i. \{y \in \text{space } M. \text{max } 0\ (s\ i\ y) \neq 0\} \subseteq \{y \in \text{space } M. s\ i\ y \neq 0\}$

unfolding *max-def* **by** *force*

moreover **have** $\wedge i. \{y \in \text{space } M. s\ i\ y \neq 0\} \in \text{sets } M$ **using** $*$ **by** *measurable*

ultimately **have** $\wedge i. \text{emeasure } M\ \{y \in \text{space } M. \text{max } 0\ (s\ i\ y) \neq 0\} \leq$

$\text{emeasure } M\ \{y \in \text{space } M. s\ i\ y \neq 0\}$ **by** (*rule emeasure-mono*)

hence $**:\wedge i. \text{emeasure } M\ \{y \in \text{space } M. \text{max } 0\ (s\ i\ y) \neq 0\} \neq \infty$ **using** $*(2)$

by (*auto intro: order.strict-trans1 simp add: top.not-eq-extremum*)

have $\wedge x. x \in \text{space } M \implies (\lambda i. \text{max } 0\ (s\ i\ x)) \longrightarrow \text{max } 0\ (f\ x)$ **using** $*(3)$

tendsto-max **by** *blast*

moreover **have** $\wedge x\ i. x \in \text{space } M \implies \text{norm } (\text{max } 0\ (s\ i\ x)) \leq \text{norm } (2 *_R$

$f\ x)$ **using** $*(4)$ **unfolding** *max-def* **by** *auto*

ultimately **have** *tendsto*: $(\lambda i. \text{integral}^L\ M\ (\lambda x. \text{max } 0\ (s\ i\ x))) \longrightarrow \text{integral}^L\ M\ (\lambda x. \text{max } 0\ (f\ x))$

using *borel-measurable-simple-function simple integrable* **by** (*intro integral-dominated-convergence*[*OF max(1) - integrable-norm*][*OF integrable-scaleR-right*], *of - 2 f*], *blast+*)

$\{$

fix $h :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assume *simple-function* $M\ h$ $\text{emeasure } M\ \{y \in \text{space } M. h\ y \neq 0\} \neq \infty \wedge x.$

$x \in \text{space } M \longrightarrow h\ x \geq 0$

hence $*$: $\text{integral}^L\ M\ h \geq 0$

proof (*induct rule: integrable-simple-function-induct-nn*)

case (*cong f g*)

then show *?case* **using** *Bochner-Integration.integral-cong* **by** *force*

next

case (*indicator A y*)

hence $A \neq \{\} \implies y \geq 0$ **using** *sets.sets-into-space* **by** *fastforce*

then show *?case* **using** *indicator* **by** (*cases A = {}*), *auto simp add: scaleR-nonneg-nonneg*)

next

case (*add f g*)


```

    then show ?case by (simp add: integrable-simple-function)
  qed
}
thus ?thesis using LIMSEQ-le-const[OF tendsto, of 0] ** simple by fastforce
qed
also have ... = integralL M f using nonneg by (auto intro: integral-cong-AE)
finally show ?thesis .
qed (simp add: not-integrable-integral-eq)

```

lemma *integral-mono-AE-banach*:

```

  fixes f g :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
    dered-real-vector}
  assumes integrable M f integrable M g AE x in M. f x ≤ g x
  shows integralL M f ≤ integralL M g
  using integral-nonneg-AE-banach[of λx. g x - f x] assms Bochner-Integration.integral-diff[OF
    assms(1,2)] by force

```

lemma *integral-mono-banach*:

```

  fixes f g :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
    dered-real-vector}
  assumes integrable M f integrable M g ∧ x. x ∈ space M ⇒ f x ≤ g x
  shows integralL M f ≤ integralL M g
  using integral-mono-AE-banach assms by blast

```

5 Integrability and Measurability of the Diameter

context

```

  fixes s :: nat ⇒ 'a ⇒ 'b :: {second-countable-topology, banach} and M
  assumes bounded: ∧x. x ∈ space M ⇒ bounded (range (λi. s i x))
begin

```

lemma *borel-measurable-diameter*:

```

  assumes [measurable]: ∧i. (s i) ∈ borel-measurable M
  shows (λx. diameter {s i x | i. n ≤ i}) ∈ borel-measurable M

```

proof –

```

  have {s i x | i. N ≤ i} ≠ {} for x N by blast
  hence diameter-SUP: diameter {s i x | i. N ≤ i} = (SUP (i, j) ∈ {N..} × {N..}.
    dist (s i x) (s j x)) for x N unfolding diameter-def by (auto intro!: arg-cong[of -
    - Sup])

```

```

  have case-prod dist '({s i x | i. n ≤ i} × {s i x | i. n ≤ i}) = ((λ(i, j). dist (s i
    x) (s j x)) '({n..} × {n..})) for x by fast

```

```

  hence *: (λx. diameter {s i x | i. n ≤ i}) = (λx. Sup ((λ(i, j). dist (s i x) (s j
    x)) '({n..} × {n..}))) using diameter-SUP by (simp add: case-prod-beta')

```

```

  have bounded ((λ(i, j). dist (s i x) (s j x)) '({n..} × {n..})) if x ∈ space M for
    x by (rule bounded-imp-dist-bounded[OF bounded, OF that])

```

```

  hence bdd: bdd-above ((λ(i, j). dist (s i x) (s j x)) '({n..} × {n..})) if x ∈ space
    M for x using that bounded-imp-bdd-above by presburger

```

have $\text{fst } p \in \text{borel-measurable } M$ **and** $p \in \text{borel-measurable } M$ **if** $p \in s \text{ ' } \{n..\} \times s \text{ ' } \{n..\}$ **for** p **using** *that* **by** *fastforce+*
hence $(\lambda x. \text{fst } p \ x - \text{snd } p \ x) \in \text{borel-measurable } M$ **if** $p \in s \text{ ' } \{n..\} \times s \text{ ' } \{n..\}$ **for** p **using** *that* **borel-measurable-diff** **by** *simp*
hence $(\lambda x. \text{case } p \text{ of } (f, g) \Rightarrow \text{dist } (f \ x) (g \ x)) \in \text{borel-measurable } M$ **if** $p \in s \text{ ' } \{n..\} \times s \text{ ' } \{n..\}$ **for** p **unfolding** *dist-norm* **using** *that* **by** *measurable*
moreover **have** *countable* $(s \text{ ' } \{n..\} \times s \text{ ' } \{n..\})$ **by** *(intro countable-SIGMA countable-image, auto)*
ultimately show *?thesis* **unfolding** ***** **by** *(auto intro!: borel-measurable-cSUP bdd)*
qed

lemma *integrable-bound-diameter*:

fixes $f :: 'a \Rightarrow \text{real}$
assumes *integrable* $M \ f$
and *[measurable]*: $\bigwedge i. (s \ i) \in \text{borel-measurable } M$
and $\bigwedge x \ i. x \in \text{space } M \implies \text{norm } (s \ i \ x) \leq f \ x$
shows *integrable* $M \ (\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\})$
proof –
have $\{s \ i \ x \mid i. N \leq i\} \neq \{\}$ **for** $x \ N$ **by** *blast*
hence *diameter-SUP*: $\text{diameter } \{s \ i \ x \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s \ i \ x) (s \ j \ x))$ **for** $x \ N$ **unfolding** *diameter-def* **by** *(auto intro!: arg-cong[of - - Sup])*
{
fix x **assume** $x \in \text{space } M$
let $?S = (\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) \text{ ' } (\{n..\} \times \{n..\})$
have *case-prod* $\text{dist ' } (\{s \ i \ x \mid i. n \leq i\} \times \{s \ i \ x \mid i. n \leq i\}) = (\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) \text{ ' } (\{n..\} \times \{n..\})$ **by** *fast*
hence $*$: $\text{diameter } \{s \ i \ x \mid i. n \leq i\} = \text{Sup } ?S$ **using** *diameter-SUP* **by** *(simp add: case-prod-beta')*

have *bounded* $?S$ **by** *(rule bounded-imp-dist-bounded[OF bounded[OF x]])*
hence $\text{Sup } S\text{-nonneg:} 0 \leq \text{Sup } ?S$ **by** *(auto intro!: cSup-upper2 x bounded-imp-bdd-above)*

have $\text{dist } (s \ i \ x) (s \ j \ x) \leq 2 * f \ x$ **for** $i \ j$ **by** *(intro dist-triangle2[THEN order-trans, of - 0]) (metis norm-conv-dist assms(3) x add-mono mult-2)*
hence $\forall c \in ?S. c \leq 2 * f \ x$ **by** *force*
hence $\text{Sup } ?S \leq 2 * f \ x$ **by** *(intro cSup-least, auto)*
hence $\text{norm } (\text{Sup } ?S) \leq 2 * \text{norm } (f \ x)$ **using** *Sup-S-nonneg* **by** *auto*
also have $\dots = \text{norm } (2 *_{\mathbb{R}} f \ x)$ **by** *simp*
finally have $\text{norm } (\text{diameter } \{s \ i \ x \mid i. n \leq i\}) \leq \text{norm } (2 *_{\mathbb{R}} f \ x)$ **unfolding**
 $*$.
}
hence $\text{AE } x \text{ in } M. \text{norm } (\text{diameter } \{s \ i \ x \mid i. n \leq i\}) \leq \text{norm } (2 *_{\mathbb{R}} f \ x)$ **by** *blast*
thus *integrable* $M \ (\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\})$ **using** *borel-measurable-diameter*
by *(intro Bochner-Integration.integrable-bound[OF assms(1)[THEN integrable-scaleR-right[of 2]]], measurable)*
qed
end

6 Auxiliary Lemmas for Integrals on a Set

lemma *set-integral-scaleR-left*:

assumes $A \in \text{sets } M \text{ } c \neq 0 \implies \text{integrable } M f$
shows $\text{LINT } t:A|M. f t *_{\mathbb{R}} c = (\text{LINT } t:A|M. f t) *_{\mathbb{R}} c$
unfolding *set-lebesgue-integral-def*
using *integrable-mult-indicator[OF assms]*
by (*subst integral-scaleR-left[symmetric], auto*)

lemma *nn-set-integral-eq-set-integral*:

assumes [*measurable*]:*integrable* $M f$
and $\forall x \in A \text{ in } M. 0 \leq f x \wedge A \in \text{sets } M$
shows $(\int^{+x \in A. f x} \partial M) = (\int x \in A. f x \partial M)$
proof –
have $(\int^{+x. \text{indicator } A x} *_{\mathbb{R}} f x \partial M) = (\int x \in A. f x \partial M)$
unfolding *set-lebesgue-integral-def* **using** *assms(2)* **by** (*intro nn-integral-eq-integral[of*
*- \lambda x. indicat-real A x *_{\mathbb{R}} f x]*, *blast intro: assms integrable-mult-indicator, fastforce*)
moreover have $(\int^{+x. \text{indicator } A x} *_{\mathbb{R}} f x \partial M) = (\int^{+x \in A. f x} \partial M)$ **by** (*metis*
ennreal-0 indicator-simps(1) indicator-simps(2) mult.commute mult-1 mult-zero-left
real-scaleR-def)
ultimately show *?thesis* **by** *argo*
qed

lemma *set-integral-restrict-space*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{banach, second-countable-topology}\}$
assumes $\Omega \cap \text{space } M \in \text{sets } M$
shows $\text{set-lebesgue-integral } (\text{restrict-space } M \Omega) A f = \text{set-lebesgue-integral } M A$
 $(\lambda x. \text{indicator } \Omega x *_{\mathbb{R}} f x)$
unfolding *set-lebesgue-integral-def*
by (*subst integral-restrict-space, auto intro!: integrable-mult-indicator assms simp:*
mult.commute)

lemma *set-integral-const*:

fixes $c :: 'b :: \{\text{banach, second-countable-topology}\}$
assumes $A \in \text{sets } M \text{ } \text{emeasure } M A \neq \infty$
shows $\text{set-lebesgue-integral } M A (\lambda \cdot. c) = \text{measure } M A *_{\mathbb{R}} c$
unfolding *set-lebesgue-integral-def*
using *assms* **by** (*metis has-bochner-integral-indicator has-bochner-integral-integral-eq*
infinity-ennreal-def less-top)

lemma *set-integral-mono-banach*:

fixes $f g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$
dered-real-vector\}
assumes *set-integrable* $M A f$ *set-integrable* $M A g$
 $\bigwedge x. x \in A \implies f x \leq g x$
shows $(\text{LINT } x:A|M. f x) \leq (\text{LINT } x:A|M. g x)$
using *assms* **unfolding** *set-integrable-def set-lebesgue-integral-def*
by (*auto intro: integral-mono-banach split: split-indicator*)

lemma *set-integral-mono-AE-banach*:

fixes $f\ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes *set-integrable* $M\ A\ f$ *set-integrable* $M\ A\ g$ $\text{AE } x \in A \text{ in } M. f\ x \leq g\ x$

shows *set-lebesgue-integral* $M\ A\ f \leq \text{set-lebesgue-integral } M\ A\ g$ **using** *assms*
unfolding *set-lebesgue-integral-def* **by** (*auto simp add: set-integrable-def intro!:*
integral-mono-AE-banach[*of* $M\ \lambda x. \text{indicator } A\ x *_R f\ x\ \lambda x. \text{indicator } A\ x *_R g\ x$],
simp add: indicator-def)

7 Averaging Theorem

lemma *balls-countable-basis*:

obtains $D :: 'a :: \{\text{metric-space, second-countable-topology}\}$ *set*

where *topological-basis* (*case-prod ball* ' ($D \times (\mathbb{Q} \cap \{0 < ..\})$)))

and *countable* D

and $D \neq \{\}$

proof –

obtain $D :: 'a$ *set* **where** *dense-subset: countable* $D\ D \neq \{\}$ [*open* $U; U \neq \{\}$]

$\Rightarrow \exists y \in D. y \in U$ **for** U **using** *countable-dense-exists* **by** *blast*

have *topological-basis* (*case-prod ball* ' ($D \times (\mathbb{Q} \cap \{0 < ..\})$)))

proof (*intro topological-basis-iff*[*THEN iffD2*], *fast, clarify*)

fix U **and** $x :: 'a$ **assume** *asm: open* $U\ x \in U$

obtain e **where** $e: e > 0$ *ball* $x\ e \subseteq U$ **using** *asm openE* **by** *blast*

obtain y **where** $y: y \in D\ y \in \text{ball } x\ (e / 3)$ **using** *dense-subset*(\mathcal{B})[*OF open-ball,*
of $x\ e / 3$] *centre-in-ball*[*THEN iffD2, OF divide-pos-pos*[*OF e(1), of 3*]] **by** *force*

obtain r **where** $r: r \in \mathbb{Q} \cap \{e/3 < .. < e/2\}$ **unfolding** *Rats-def* **using** *of-rat-dense*[*OF*
divide-strict-left-mono[*OF - e(1)*], *of 2 3*] **by** *auto*

have $*$: $x \in \text{ball } y\ r$ **using** $r\ y$ **by** (*simp add: dist-commute*)

hence $\text{ball } y\ r \subseteq U$ **using** r **by** (*intro order-trans*[*OF - e(2)*], *simp, metric*)

moreover **have** $\text{ball } y\ r \in (\text{case-prod ball ' } (D \times (\mathbb{Q} \cap \{0 < ..\})))$ **using** $y(1)$

r **by** *force*

ultimately show $\exists B' \in (\text{case-prod ball ' } (D \times (\mathbb{Q} \cap \{0 < ..\}))). x \in B' \wedge B' \subseteq U$
using $*$ **by** *meson*

qed

thus *?thesis* **using** *that dense-subset* **by** *blast*

qed

context *sigma-finite-measure*

begin

lemma *sigma-finite-measure-induct*[*case-names finite-measure, consumes 0*]:

assumes $\bigwedge(N :: 'a \text{ measure})\ \Omega. \text{finite-measure } N$

$\Rightarrow N = \text{restrict-space } M\ \Omega$

$\Rightarrow \Omega \in \text{sets } M$

$\Rightarrow \text{emeasure } N\ \Omega \neq \infty$

$\Rightarrow \text{emeasure } N\ \Omega \neq 0$

$\Rightarrow \text{almost-everywhere } N\ Q$

and [*measurable*]: *Measurable.pred* $M\ Q$

shows *almost-everywhere* $M \ Q$
proof –
have $*$: *almost-everywhere* $N \ Q$ **if** *finite-measure* $N \ N = \text{restrict-space } M \ \Omega \ \Omega$
 $\in \text{sets } M \ \text{emeasure } N \ \Omega \neq \infty$ **for** $N \ \Omega$ **using** *that* **by** (*cases* *emeasure* $N \ \Omega = 0$,
auto *intro*: *emeasure-0-AE* *assms*(1))

obtain $A :: \text{nat} \Rightarrow 'a \ \text{set}$ **where** A : *range* $A \subseteq \text{sets } M \ (\bigcup i. A \ i) = \text{space } M$ **and**
emeasure-finite: *emeasure* $M \ (A \ i) \neq \infty$ **for** i **using** *sigma-finite* **by** *metis*
note $A(1)$ [*measurable*]
have *space-restr*: *space* (*restrict-space* $M \ (A \ i)$) = $A \ i$ **for** i **unfolding** *space-restrict-space*
by *simp*
{
 fix i
 have $*$: $\{x \in A \ i \cap \text{space } M. \ Q \ x\} = \{x \in \text{space } M. \ Q \ x\} \cap (A \ i)$ **by** *fast*
 have *Measurable.pred* (*restrict-space* $M \ (A \ i)$) Q **using** A **by** (*intro* *measurableI*,
auto *simp* *add*: *space-restr* *intro*!: *sets-restrict-space-iff*[*THEN* *iffD2*], *measurable*,
auto)
}
note *this*[*measurable*]
{
 fix i
 have *finite-measure* (*restrict-space* $M \ (A \ i)$) **using** *emeasure-finite* **by** (*intro*
finite-measureI, *subst* *space-restr*, *subst* *emeasure-restrict-space*, *auto*)
 hence *emeasure* (*restrict-space* $M \ (A \ i)$) $\{x \in A \ i. \neg Q \ x\} = 0$ **using** *emea-*
sure-finite **by** (*intro* *AE-iff-measurable*[*THEN* *iffD1*, *OF* - - *], *measurable*, *subst*
space-restr[*symmetric*], *intro* *sets.top*, *auto* *simp* *add*: *emeasure-restrict-space*)
 hence *emeasure* $M \ \{x \in A \ i. \neg Q \ x\} = 0$ **by** (*subst* *emeasure-restrict-space*[*symmetric*],
auto)
}
hence *emeasure* $M \ (\bigcup i. \{x \in A \ i. \neg Q \ x\}) = 0$ **by** (*intro* *emeasure-UN-eq-0*,
auto)
moreover **have** $(\bigcup i. \{x \in A \ i. \neg Q \ x\}) = \{x \in \text{space } M. \neg Q \ x\}$ **using** A **by**
auto
ultimately show *?thesis* **by** (*intro* *AE-iff-measurable*[*THEN* *iffD2*], *auto*)
qed

lemma *averaging-theorem*:
fixes $f :: \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes [*measurable*]: *integrable* $M \ f$
and *closed*: *closed* S
and $\bigwedge A. A \in \text{sets } M \implies \text{measure } M \ A > 0 \implies (1 / \text{measure } M \ A) *_R$
set-lebesgue-integral $M \ A \ f \in S$
shows *AE* x *in* $M. f \ x \in S$
proof (*induct* *rule*: *sigma-finite-measure-induct*)
case (*finite-measure* $N \ \Omega$)

interpret *finite-measure* N **by** (*rule* *finite-measure*)

have *integrable[measurable]: integrable N f* **using** *assms finite-measure* **by** (*auto simp: integrable-restrict-space integrable-mult-indicator*)
have *average: (1 / Sigma-Algebra.measure N A) *_R set-lebesgue-integral N A f* $\in S$ **if** $A \in \text{sets } N \text{ measure } N A > 0$ **for** A
proof –
have $*$: $A \in \text{sets } M$ **using** *that* **by** (*simp add: sets-restrict-space-iff finite-measure*)
have $A = A \cap \Omega$ **by** (*metis finite-measure(2,3) inf.orderE sets.sets-into-space space-restrict-space that(1)*)
hence *set-lebesgue-integral N A f = set-lebesgue-integral M A f* **unfolding** *finite-measure* **by** (*subst set-integral-restrict-space, auto simp add: finite-measure set-lebesgue-integral-def indicator-inter-arith[symmetric]*)
moreover **have** *measure N A = measure M A* **using** *that* **by** (*auto intro!: measure-restrict-space simp add: finite-measure sets-restrict-space-iff*)
ultimately show *?thesis* **using** *that * assms(3)* **by** *presburger*
qed

obtain $D :: 'b \text{ set}$ **where** *balls-basis: topological-basis (case-prod ball ‘ (D × (Q ∩ {0<..})))* **and** *countable-D: countable D* **using** *balls-countable-basis* **by** *blast*
have *countable-balls: countable (case-prod ball ‘ (D × (Q ∩ {0<..})))* **using** *countable-rat countable-D* **by** *blast*

obtain B **where** *B-balls: B ⊆ case-prod ball ‘ (D × (Q ∩ {0<..})) ∪ B = −S* **using** *topological-basis[THEN iffD1, OF balls-basis] open-Compl[OF assms(2)]* **by** *meson*
hence *countable-B: countable B* **using** *countable-balls countable-subset* **by** *fast*

define b **where** $b = \text{from-nat-into } (B \cup \{\{\}\})$
have $B \cup \{\{\}\} \neq \{\}$ **by** *simp*
have *range-b: range b = B ∪ {\{\}}* **using** *countable-B* **by** (*auto simp add: b-def intro!: range-from-nat-into*)
have *open-b: open (b i) for i* **unfolding** *b-def* **using** *B-balls open-ball from-nat-into[of B ∪ {\{\}} i]* **by** *force*
have *Union-range-b: ∪(range b) = −S* **using** *B-balls range-b* **by** *simp*

{
fix $v \ r$ **assume** *ball-in-Compl: ball v r ⊆ −S*
define A **where** $A = f - ' \text{ball } v \ r \cap \text{space } N$
have *dist-less: dist (f x) v < r* **if** $x \in A$ **for** x **using** *that* **unfolding** *A-def vimage-def* **by** (*simp add: dist-commute*)
hence *AE-less: AE x ∈ A in N. norm (f x − v) < r* **by** (*auto simp add: dist-norm*)
have $*$: $A \in \text{sets } N$ **unfolding** *A-def* **by** *simp*
have *emeasure N A = 0*
proof –
{
assume *asm: emeasure N A > 0*
hence *measure-pos: measure N A > 0* **unfolding** *emeasure-eq-measure* **by** *simp*
hence $(1 / \text{measure } N A) *_{\text{R}} \text{set-lebesgue-integral } N A f - v = (1 / \text{measure } N$

$A) *_{\mathcal{R}} \text{set-lebesgue-integral } N \ A \ (\lambda x. f \ x - v) \text{ using } \text{integrable integrable-const} * \text{ by}$
 $(\text{subst set-integral-diff}(2), \text{auto simp add: set-integrable-def set-integral-const}[OF *])$
 $\text{algebra-simps intro!: integrable-mult-indicator})$
 $\text{moreover have norm } (\int_{x \in A. (f \ x - v) \partial N) \leq (\int_{x \in A. \text{norm } (f \ x - v) \partial N) \text{ using } * \text{ by } (\text{auto intro!: integral-norm-bound}[of \ N \ \lambda x. \text{indicator } A \ x$
 $*_{\mathcal{R}} (f \ x - v), \text{ THEN order-trans}] \text{ integrable-mult-indicator integrable simp add:}$
 $\text{set-lebesgue-integral-def})$
 $\text{ultimately have norm } ((1 / \text{measure } N \ A) *_{\mathcal{R}} \text{set-lebesgue-integral } N \ A \ f$
 $- v) \leq \text{set-lebesgue-integral } N \ A \ (\lambda x. \text{norm } (f \ x - v)) / \text{measure } N \ A \text{ using } \text{asm}$
 $\text{by } (\text{auto intro: divide-right-mono})$
 $\text{also have } \dots < \text{set-lebesgue-integral } N \ A \ (\lambda x. r) / \text{measure } N \ A$
 $\text{unfolding set-lebesgue-integral-def}$
 $\text{using } \text{asm} * \text{integrable integrable-const AE-less measure-pos}$
 $\text{by } (\text{intro divide-strict-right-mono integral-less-AE}[of \ - \ A] \text{ integrable-mult-indicator})$
 $(\text{fastforce simp add: dist-less dist-norm indicator-def})+$
 $\text{also have } \dots = r \text{ using } * \text{measure-pos by } (\text{simp add: set-integral-const})$
 $\text{finally have dist } ((1 / \text{measure } N \ A) *_{\mathcal{R}} \text{set-lebesgue-integral } N \ A \ f) \ v < r$
 $\text{by } (\text{subst dist-norm})$
 $\text{hence False using average}[OF * \text{measure-pos}] \text{ by } (\text{metis ComplD dist-commute}$
 $\text{in-mono mem-ball ball-in-Compl})$
 $\}$
 $\text{thus ?thesis by fastforce}$
 qed
 $\}$
 $\text{note } * = \text{this}$
 $\{$
 $\text{fix } b' \text{ assume } b' \in B$
 $\text{hence ball-subset-Compl: } b' \subseteq -S \text{ and ball-radius-pos: } \exists v \in D. \exists r > 0. b' =$
 $\text{ball } v \ r \text{ using B-balls by } (\text{blast, fast})$
 $\}$
 $\text{note } ** = \text{this}$
 $\text{hence } \text{emeasure } N \ (f - ' b \ i \cap \text{space } N) = 0 \text{ for } i \text{ by } (\text{cases } b \ i = \{\}, \text{ simp})$
 $(\text{metis UnE singletonD } * \text{ range-b}[THEN \text{eq-refl}, \text{ THEN range-subsetD}])$
 $\text{hence } \text{emeasure } N \ (\bigcup i. f - ' b \ i \cap \text{space } N) = 0 \text{ using open-b by } (\text{intro}$
 $\text{emeasure-UN-eq-0}) \text{ fastforce+}$
 $\text{moreover have } (\bigcup i. f - ' b \ i \cap \text{space } N) = f - ' (\bigcup (\text{range } b)) \cap \text{space } N \text{ by}$
 blast
 $\text{ultimately have } \text{emeasure } N \ (f - ' (-S) \cap \text{space } N) = 0 \text{ using Union-range-b}$
 by argo
 $\text{hence } \text{AE } x \text{ in } N. f \ x \notin -S \text{ using open-Compl}[OF \text{assms}(2)] \text{ by } (\text{intro AE-iff-measurable}[THEN$
 $\text{iffD2}], \text{ auto})$
 $\text{thus ?case by force}$
 $\text{qed } (\text{simp add: pred-sets2}[OF \text{borel-closed}] \text{assms}(2))$

lemma density-zero:
 $\text{fixes } f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
 $\text{assumes integrable } M \ f$
 $\text{and density-0: } \bigwedge A. A \in \text{sets } M \implies \text{set-lebesgue-integral } M \ A \ f = 0$
 $\text{shows AE } x \text{ in } M. f \ x = 0$

```

using averaging-theorem[OF assms(1), of {0}] assms(2)
by (simp add: scaleR-nonneg-nonneg)

lemma density-unique-banach:
  fixes f f'::'a  $\Rightarrow$  'b::{second-countable-topology, banach}
  assumes integrable M f integrable M f'
  and density-eq:  $\bigwedge A. A \in \text{sets } M \implies \text{set-lebesgue-integral } M A f = \text{set-lebesgue-integral } M A f'$ 
  shows AE x in M. f x = f' x
proof -
  {
    fix A assume asm: A  $\in$  sets M
    hence LINT x|M. indicat-real A x *R (f x - f' x) = 0 using density-eq
    assms(1,2) by (simp add: set-lebesgue-integral-def algebra-simps Bochner-Integration.integral-diff[OF integrable-mult-indicator(1,1)])
  }
  thus ?thesis using density-zero[OF Bochner-Integration.integrable-diff[OF assms(1,2)]]
by (simp add: set-lebesgue-integral-def)
qed

lemma density-nonneg:
  fixes f::-  $\Rightarrow$  'b::{second-countable-topology, banach, linorder-topology, ordered-real-vector}
  assumes integrable M f
  and  $\bigwedge A. A \in \text{sets } M \implies \text{set-lebesgue-integral } M A f \geq 0$ 
  shows AE x in M. f x  $\geq$  0
  using averaging-theorem[OF assms(1), of {0..}, OF closed-atLeast] assms(2)
  by (simp add: scaleR-nonneg-nonneg)

corollary integral-nonneg-AE-eq-0-iff-AE:
  fixes f :: 'a  $\Rightarrow$  'b :: {second-countable-topology, banach, linorder-topology, ordered-real-vector}
  assumes f[measurable]: integrable M f and nonneg: AE x in M. 0  $\leq$  f x
  shows integralL M f = 0  $\longleftrightarrow$  (AE x in M. f x = 0)
proof
  assume *: integralL M f = 0
  {
    fix A assume asm: A  $\in$  sets M
    have 0  $\leq$  integralL M ( $\lambda x$ . indicator A x *R f x) using nonneg by (subst integral-zero[of M, symmetric], intro integral-mono-AE-banach integrable-mult-indicator asm f integrable-zero, auto simp add: indicator-def)
    moreover have ...  $\leq$  integralL M f using nonneg by (intro integral-mono-AE-banach integrable-mult-indicator asm f, auto simp add: indicator-def)
    ultimately have set-lebesgue-integral M A f = 0 unfolding set-lebesgue-integral-def using * by force
  }
  thus AE x in M. f x = 0 by (intro density-zero f, blast)
qed (auto simp add: integral-eq-zero-AE)

corollary integral-eq-mono-AE-eq-AE:

```



```

fixes  $f\ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$ 
 $\text{dered-real-vector}\}$ 
assumes  $\text{integrable } M\ f\ \text{integrable } M\ g\ \text{integral}^L\ M\ f = \text{integral}^L\ M\ g\ AE\ x\ \text{in}$ 
 $M. f\ x \leq g\ x$ 
shows  $AE\ x\ \text{in } M. f\ x = g\ x$ 
proof –
  define  $h$  where  $h = (\lambda x. g\ x - f\ x)$ 
  have  $AE\ x\ \text{in } M. h\ x = 0$  unfolding  $h\text{-def}$  using  $assms$  by  $(subst\ \text{integral-nonneg-AE-eq-0-iff-AE[symmetric]})\ auto$ 
  then show  $?thesis$  unfolding  $h\text{-def}$  by  $auto$ 
qed

end

end

```

```

theory Conditional-Expectation-Banach
imports HOL-Probability.Conditional-Expectation HOL-Probability.Independent-Family
Bochner-Integration-Supplement
begin

```

8 Conditional Expectation in Banach Spaces

```

definition  $has\text{-}cond\text{-}exp :: 'a\ measure \Rightarrow 'a\ measure \Rightarrow ('a \Rightarrow 'b :: \{\text{real-normed-vector,}$ 
 $\text{second-countable-topology}\}) \Rightarrow bool$  where
   $has\text{-}cond\text{-}exp\ M\ F\ f\ g = ((\forall A \in sets\ F. (\int x \in A. f\ x\ \partial M) = (\int x \in A. g\ x\ \partial M))$ 
 $\wedge\ integrable\ M\ f$ 
 $\wedge\ integrable\ M\ g$ 
 $\wedge\ g \in \text{borel-measurable } F)$ 

```

```

lemma  $has\text{-}cond\text{-}expI'$ :
assumes  $\bigwedge A. A \in sets\ F \implies (\int x \in A. f\ x\ \partial M) = (\int x \in A. g\ x\ \partial M)$ 
 $integrable\ M\ f$ 
 $integrable\ M\ g$ 
 $g \in \text{borel-measurable } F$ 
shows  $has\text{-}cond\text{-}exp\ M\ F\ f\ g$ 
using  $assms$  unfolding  $has\text{-}cond\text{-}exp\text{-}def$  by  $simp$ 

```

```

lemma  $has\text{-}cond\text{-}expD$ :
assumes  $has\text{-}cond\text{-}exp\ M\ F\ f\ g$ 
shows  $\bigwedge A. A \in sets\ F \implies (\int x \in A. f\ x\ \partial M) = (\int x \in A. g\ x\ \partial M)$ 
 $integrable\ M\ f$ 
 $integrable\ M\ g$ 
 $g \in \text{borel-measurable } F$ 
using  $assms$  unfolding  $has\text{-}cond\text{-}exp\text{-}def$  by  $simp+$ 

```

definition *cond-exp* :: 'a measure \Rightarrow 'a measure \Rightarrow ('a \Rightarrow 'b) \Rightarrow ('a \Rightarrow 'b::{banach, second-countable-topology}) **where**
cond-exp *M F f* = (if $\exists g. \text{has-cond-exp } M F f g$ then (SOME *g. has-cond-exp M F f g*) else ($\lambda-. 0$))

lemma *borel-measurable-cond-exp*[*measurable*]: *cond-exp M F f* \in *borel-measurable F*
by (metis *cond-exp-def someI has-cond-exp-def borel-measurable-const*)

lemma *integrable-cond-exp*[*intro*]: *integrable M (cond-exp M F f)*
by (metis *cond-exp-def has-cond-expD(3) integrable-zero someI*)

lemma *set-integrable-cond-exp*[*intro*]:
assumes *A* \in *sets M*
shows *set-integrable M A (cond-exp M F f)* **using** *integrable-mult-indicator[OF assms integrable-cond-exp, of F f]* **by** (auto simp add: *set-integrable-def intro!*: *integrable-mult-indicator[OF assms integrable-cond-exp]*)

lemma *has-cond-exp-self*:
assumes *integrable M f*
shows *has-cond-exp M (vimage-algebra (space M) f borel) f f*
using *assms* **by** (auto *intro!*: *has-cond-expI' measurable-vimage-algebra1*)

lemma *has-cond-exp-sets-cong*:
assumes *sets F = sets G*
shows *has-cond-exp M F = has-cond-exp M G*
using *assms* **unfolding** *has-cond-exp-def* **by** *force*

lemma *cond-exp-sets-cong*:
assumes *sets F = sets G*
shows $\forall x \text{ in } M. \text{cond-exp } M F f x = \text{cond-exp } M G f x$
by (*intro AE-I2, simp add: cond-exp-def has-cond-exp-sets-cong[OF assms, of M]*)

context *sigma-finite-subalgebra*
begin

lemma *borel-measurable-cond-exp'*[*measurable*]: *cond-exp M F f* \in *borel-measurable M*
by (metis *cond-exp-def someI has-cond-exp-def borel-measurable-const subalg measurable-from-subalg*)

lemma *cond-exp-null*:
assumes $\nexists g. \text{has-cond-exp } M F f g$
shows *cond-exp M F f* = ($\lambda-. 0$)
unfolding *cond-exp-def* **using** *assms* **by** *argo*

lemma *has-cond-exp-nested-subalg*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes $\text{subalgebra } G \ F \ \text{has-cond-exp } M \ F \ f \ h \ \text{has-cond-exp } M \ G \ f \ h'$
shows $\text{has-cond-exp } M \ F \ h' \ h$
by $(\text{intro } \text{has-cond-expI}') (\text{metis } \text{assms } \text{has-cond-expD } \text{in-mono } \text{subalgebra-def}) +$

lemma $\text{has-cond-exp-charact}$:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes $\text{has-cond-exp } M \ F \ f \ g$
shows $\text{has-cond-exp } M \ F \ f \ (\text{cond-exp } M \ F \ f)$
 $AE \ x \ \text{in } M. \ \text{cond-exp } M \ F \ f \ x = g \ x$

proof –
show $\text{cond-exp: } \text{has-cond-exp } M \ F \ f \ (\text{cond-exp } M \ F \ f)$ **using** $\text{assms } \text{someI}$
 cond-exp-def **by** metis
let $?MF = \text{restr-to-subalg } M \ F$
interpret $\text{sigma-finite-measure } ?MF$ **by** $(\text{rule } \text{sigma-fin-subalg})$
{
fix A **assume** $A \in \text{sets } ?MF$
then have $[\text{measurable}]: A \in \text{sets } F$ **using** $\text{sets-restr-to-subalg}[OF \ \text{subalg}]$ **by**
 simp
have $(\int x \in A. g \ x \ \partial ?MF) = (\int x \in A. g \ x \ \partial M)$ **using** $\text{assms } \text{subalg}$ **by** $(\text{auto } \text{simp } \text{add: } \text{integral-subalgebra2 } \text{set-lebesgue-integral-def } \text{dest!}: \text{has-cond-expD})$
also have $\dots = (\int x \in A. \text{cond-exp } M \ F \ f \ x \ \partial M)$ **using** $\text{assms } \text{cond-exp}$ **by**
 $(\text{simp } \text{add: } \text{has-cond-exp-def})$
also have $\dots = (\int x \in A. \text{cond-exp } M \ F \ f \ x \ \partial ?MF)$ **using** subalg **by** $(\text{auto } \text{simp } \text{add: } \text{integral-subalgebra2 } \text{set-lebesgue-integral-def})$
finally have $(\int x \in A. g \ x \ \partial ?MF) = (\int x \in A. \text{cond-exp } M \ F \ f \ x \ \partial ?MF)$ **by**
 simp
}
hence $AE \ x \ \text{in } ?MF. \ \text{cond-exp } M \ F \ f \ x = g \ x$ **using** $\text{cond-exp } \text{assms } \text{subalg}$ **by**
 $(\text{intro } \text{density-unique-banach}, \text{auto } \text{dest: } \text{has-cond-expD } \text{intro!}: \text{integrable-in-subalg})$
then show $AE \ x \ \text{in } M. \ \text{cond-exp } M \ F \ f \ x = g \ x$ **using** $AE\text{-restr-to-subalg}[OF \ \text{subalg}]$ **by** simp
qed

lemma cond-exp-charact :
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes $\bigwedge A. A \in \text{sets } F \implies (\int x \in A. f \ x \ \partial M) = (\int x \in A. g \ x \ \partial M)$
 $\text{integrable } M \ f$
 $\text{integrable } M \ g$
 $g \in \text{borel-measurable } F$
shows $AE \ x \ \text{in } M. \ \text{cond-exp } M \ F \ f \ x = g \ x$
by $(\text{intro } \text{has-cond-exp-charact } \text{has-cond-expI}' \ \text{assms}) \ \text{auto}$

corollary $\text{cond-exp-F-meas}[\text{intro}, \text{simp}]$:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes $\text{integrable } M \ f$
 $f \in \text{borel-measurable } F$
shows $AE \ x \ \text{in } M. \ \text{cond-exp } M \ F \ f \ x = f \ x$
by $(\text{rule } \text{cond-exp-charact}, \text{auto } \text{intro: } \text{assms})$

Congruence

lemma *has-cond-exp-cong*:

assumes *integrable* $M f \bigwedge x. x \in \text{space } M \implies f x = g x$ *has-cond-exp* $M F g h$

shows *has-cond-exp* $M F f h$

proof (*intro* *has-cond-expI'*[*OF* - *assms*(1)], *goal-cases*)

case (1 *A*)

hence *set-lebesgue-integral* $M A f = \text{set-lebesgue-integral } M A g$ **by** (*intro* *set-lebesgue-integral-cong*)
(*meson* *assms*(2) *subalg* *in-mono* *subalgebra-def* *sets.sets-into-space* *subalgebra-def* *subsetD*) +

then show ?*case* **using** 1 *assms*(3) **by** (*simp* *add*: *has-cond-exp-def*)

qed (*auto* *simp* *add*: *has-cond-expD*[*OF* *assms*(3)])

lemma *cond-exp-cong*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$

assumes *integrable* $M f$ *integrable* $M g \bigwedge x. x \in \text{space } M \implies f x = g x$

shows *AE* x *in* $M. \text{cond-exp } M F f x = \text{cond-exp } M F g x$

proof (*cases* $\exists h. \text{has-cond-exp } M F f h$)

case *True*

then obtain h **where** $h: \text{has-cond-exp } M F f h \text{ has-cond-exp } M F g h$ **using** *has-cond-exp-cong* *assms* **by** *metis*

show ?*thesis* **using** $h[\text{THEN } \text{has-cond-exp-charact}(2)]$ **by** *fastforce*

next

case *False*

moreover have $\nexists h. \text{has-cond-exp } M F g h$ **using** *False* *has-cond-exp-cong* *assms* **by** *auto*

ultimately show ?*thesis* **unfolding** *cond-exp-def* **by** *auto*

qed

lemma *has-cond-exp-cong-AE*:

assumes *integrable* $M f$ *AE* x *in* $M. f x = g x$ *has-cond-exp* $M F g h$

shows *has-cond-exp* $M F f h$

using *assms*(1,2) *subalg* *subalgebra-def* *subset-iff*

by (*intro* *has-cond-expI'*, *subst* *set-lebesgue-integral-cong-AE*[*OF* - *assms*(1)[*THEN* *borel-measurable-integrable*] *borel-measurable-integrable*(1)[*OF* *has-cond-expD*(2)[*OF* *assms*(3)]]])

(*fast* *intro*: *has-cond-expD*[*OF* *assms*(3)] *integrable-cong-AE-imp*[*OF* - - *AE-symmetric*]) +

lemma *has-cond-exp-cong-AE'*:

assumes $h \in \text{borel-measurable } F$ *AE* x *in* $M. h x = h' x$ *has-cond-exp* $M F f h'$

shows *has-cond-exp* $M F f h$

using *assms*(1, 2) *subalg* *subalgebra-def* *subset-iff*

using *AE-restr-to-subalg2*[*OF* *subalg* *assms*(2)] *measurable-from-subalg*

by (*intro* *has-cond-expI'*, *subst* *set-lebesgue-integral-cong-AE*[*OF* - *measurable-from-subalg*(1,1)[*OF* *subalg*], *OF* - *assms*(1) *has-cond-expD*(4)[*OF* *assms*(3)]]])

(*fast* *intro*: *has-cond-expD*[*OF* *assms*(3)] *integrable-cong-AE-imp*[*OF* - - *AE-symmetric*]) +

lemma *cond-exp-cong-AE*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$

assumes *integrable* $M f$ *integrable* $M g$ *AE* x *in* $M. f x = g x$

shows $AE\ x\ in\ M.\ cond_exp\ M\ F\ f\ x = cond_exp\ M\ F\ g\ x$
proof (cases $\exists h.\ has_cond_exp\ M\ F\ f\ h$)
 case *True*
 then obtain h **where** $h: has_cond_exp\ M\ F\ f\ h\ has_cond_exp\ M\ F\ g\ h$ **using**
 $has_cond_exp_cong_AE\ assms$ **by** (metis (mono-tags, lifting) eventually-mono)
 show $?thesis$ **using** $h[THEN\ has_cond_exp_charact(2)]$ **by** fastforce
 next
 case *False*
 moreover have $\nexists h.\ has_cond_exp\ M\ F\ g\ h$ **using** *False* $has_cond_exp_cong_AE$
 $assms$ **by** auto
 ultimately show $?thesis$ **unfolding** $cond_exp_def$ **by** auto
qed

lemma *has-cond-exp-real*:
 fixes $f :: 'a \Rightarrow real$
 assumes *integrable* $M\ f$
 shows $has_cond_exp\ M\ F\ f\ (real_cond_exp\ M\ F\ f)$
 by (intro *has-cond-expI'*, auto intro!: *real-cond-exp-intA* $assms$)

lemma *cond-exp-real[intro]*:
 fixes $f :: 'a \Rightarrow real$
 assumes *integrable* $M\ f$
 shows $AE\ x\ in\ M.\ cond_exp\ M\ F\ f\ x = real_cond_exp\ M\ F\ f\ x$
 using *has-cond-exp-charact* *has-cond-exp-real* $assms$ **by** blast

lemma *cond-exp-cmult*:
 fixes $f :: 'a \Rightarrow real$
 assumes *integrable* $M\ f$
 shows $AE\ x\ in\ M.\ cond_exp\ M\ F\ (\lambda x.\ c * f\ x)\ x = c * cond_exp\ M\ F\ f\ x$
 using *real-cond-exp-cmult[OF* $assms(1)$, *of* $c]$ $assms(1)[THEN\ cond_exp_real]$
 $assms(1)[THEN\ integrable_mult_right,\ THEN\ cond_exp_real,\ of\ c]$ **by** fastforce

Indicator functions

lemma *has-cond-exp-indicator*:
 assumes $A \in sets\ M\ emeasure\ M\ A < \infty$
 shows $has_cond_exp\ M\ F\ (\lambda x.\ indicat_real\ A\ x *_{\mathbb{R}}\ y)\ (\lambda x.\ real_cond_exp\ M\ F\ (indicator\ A)\ x *_{\mathbb{R}}\ y)$
proof (intro *has-cond-expI'*, goal-cases)
 case (1 B)
 have $\int_{x \in B} (indicat_real\ A\ x *_{\mathbb{R}}\ y)\ \partial M = (\int_{x \in B} indicat_real\ A\ x\ \partial M) *_{\mathbb{R}}\ y$ **using** $assms$ **by** (intro *set-integral-scaleR-left*, meson 1 *in-mono* *subalg* *subalgebra-def*, blast)
 also have $\dots = \int_{x \in B} real_cond_exp\ M\ F\ (indicator\ A)\ x\ \partial M *_{\mathbb{R}}\ y$ **using** 1 $assms$ **by** (subst *real-cond-exp-intA*, auto)
 also have $\dots = \int_{x \in B} (real_cond_exp\ M\ F\ (indicator\ A)\ x *_{\mathbb{R}}\ y)\ \partial M$ **using** $assms$ **by** (intro *set-integral-scaleR-left[symmetric]*, meson 1 *in-mono* *subalg* *subalgebra-def*, blast)
 finally show $?case$.
 next

```

    case 2
    then show ?case using integrable-scaleR-left integrable-real-indicator assms by
blast
next
    case 3
    show ?case using assms by (intro integrable-scaleR-left, intro real-cond-exp-int,
blast+)
next
    case 4
    then show ?case by (intro borel-measurable-scaleR, intro Conditional-Expectation.borel-measurable-cond-exp,
simp)
qed

```

```

lemma cond-exp-indicator[intro]:
  fixes  $y :: 'b :: \{second-countable-topology, banach\}$ 
  assumes [measurable]:  $A \in sets\ M$   $emeasure\ M\ A < \infty$ 
  shows  $\forall x \in M. cond-exp\ M\ F\ (\lambda x. indicat-real\ A\ x *_{\mathbb{R}} y)\ x = cond-exp\ M\ F$ 
 $(indicator\ A)\ x *_{\mathbb{R}} y$ 
proof -
  have  $\forall x \in M. cond-exp\ M\ F\ (\lambda x. indicat-real\ A\ x *_{\mathbb{R}} y)\ x = real-cond-exp\ M\ F$ 
 $(indicator\ A)\ x *_{\mathbb{R}} y$  using has-cond-exp-indicator[OF assms] has-cond-exp-charact
by blast
  thus ?thesis using cond-exp-real[OF integrable-real-indicator, OF assms] by fast-
force
qed

```

Addition

```

lemma has-cond-exp-add:
  fixes  $f\ g :: 'a \Rightarrow 'b :: \{second-countable-topology, banach\}$ 
  assumes has-cond-exp  $M\ F\ f\ f'$  has-cond-exp  $M\ F\ g\ g'$ 
  shows has-cond-exp  $M\ F\ (\lambda x. f\ x + g\ x)\ (\lambda x. f'\ x + g'\ x)$ 
proof (intro has-cond-expI', goal-cases)
  case (1 A)
  have  $\int_{x \in A}. (f\ x + g\ x) \partial M = (\int_{x \in A}. f\ x \partial M) + (\int_{x \in A}. g\ x \partial M)$  using
assms[THEN has-cond-expD(2)] subalg 1 by (intro set-integral-add(2), auto simp
add: subalgebra-def set-integrable-def intro: integrable-mult-indicator)
  also have  $\dots = (\int_{x \in A}. f'\ x \partial M) + (\int_{x \in A}. g'\ x \partial M)$  using assms[THEN
has-cond-expD(1)[OF - 1]] by argo
  also have  $\dots = \int_{x \in A}. (f'\ x + g'\ x) \partial M$  using assms[THEN has-cond-expD(3)]
subalg 1 by (intro set-integral-add(2)[symmetric], auto simp add: subalgebra-def
set-integrable-def intro: integrable-mult-indicator)
  finally show ?case .
next
  case 2
  then show ?case by (metis Bochner-Integration.integrable-add assms has-cond-expD(2))
next
  case 3
  then show ?case by (metis Bochner-Integration.integrable-add assms has-cond-expD(3))
next

```

case 4
 then show ?case using assms borel-measurable-add has-cond-expD(4) by blast
 qed

lemma *has-cond-exp-scaleR-right*:
 fixes $f :: 'a \Rightarrow 'b::\{\text{second-countable-topology}, \text{banach}\}$
 assumes *has-cond-exp* $M F f f'$
 shows *has-cond-exp* $M F (\lambda x. c *_R f x) (\lambda x. c *_R f' x)$
 using *has-cond-expD*[*OF* *assms*] by (intro *has-cond-expI'*, auto)

lemma *cond-exp-scaleR-right*:
 fixes $f :: 'a \Rightarrow 'b::\{\text{second-countable-topology}, \text{banach}\}$
 assumes *integrable* $M f$
 shows *AE* x in M . *cond-exp* $M F (\lambda x. c *_R f x) x = c *_R \text{cond-exp } M F f x$
proof (cases $\exists f'. \text{has-cond-exp } M F f f'$)
 case True
 then show ?thesis using assms *has-cond-exp-charact* *has-cond-exp-scaleR-right*
 by metis
 next
 case False
 show ?thesis
proof (cases $c = 0$)
 case True
 then show ?thesis by simp
 next
 case *c-nonzero*: False
 have $\nexists f'. \text{has-cond-exp } M F (\lambda x. c *_R f x) f'$
proof (standard, goal-cases)
 case 1
 then obtain f' where $f': \text{has-cond-exp } M F (\lambda x. c *_R f x) f'$ by blast
 have *has-cond-exp* $M F f (\lambda x. \text{inverse } c *_R f' x)$ using *has-cond-expD*[*OF*
 f] *divideR-right*[*OF* *c-nonzero*] *assms* by (intro *has-cond-expI'*, auto)
 then show ?case using False by blast
 qed
 then show ?thesis using *cond-exp-null*[*OF* False] *cond-exp-null* by force
 qed
 qed

lemma *cond-exp-uminus*:
 fixes $f :: 'a \Rightarrow 'b::\{\text{second-countable-topology}, \text{banach}\}$
 assumes *integrable* $M f$
 shows *AE* x in M . *cond-exp* $M F (\lambda x. - f x) x = - \text{cond-exp } M F f x$
 using *cond-exp-scaleR-right*[*OF* *assms*, of -1] by force

corollary *has-cond-exp-simple*:
 fixes $f :: 'a \Rightarrow 'b::\{\text{second-countable-topology}, \text{banach}\}$
 assumes *simple-function* $M f$ *emeasure* $M \{y \in \text{space } M. f y \neq 0\} \neq \infty$
 shows *has-cond-exp* $M F f (\text{cond-exp } M F f)$
 using *assms*

proof (*induction rule: integrable-simple-function-induct*)
case (*cong f g*)
then show ?*case* **using** *has-cond-exp-cong* **by** (*metis (no-types, opaque-lifting)*
Bochner-Integration.integrable-cong has-cond-expD(2) has-cond-exp-charact(1))
next
case (*indicator A y*)
then show ?*case* **using** *has-cond-exp-charact[OF has-cond-exp-indicator]* **by** *fast*
next
case (*add u v*)
then show ?*case* **using** *has-cond-exp-add has-cond-exp-charact(1)* **by** *blast*
qed

lemma *cond-exp-contraction-real*:

fixes *f* :: 'a \Rightarrow real
assumes *integrable[measurable]: integrable M f*
shows *AE x in M. norm (cond-exp M F f x) \leq cond-exp M F ($\lambda x. \text{norm } (f x)$) x*
proof –
have *int: integrable M ($\lambda x. \text{norm } (f x)$)* **using** *assms* **by** *blast*
have *: *AE x in M. 0 \leq cond-exp M F ($\lambda x. \text{norm } (f x)$) x* **using** *cond-exp-real[THEN*
AE-symmetric, OF integrable-norm[OF integrable]] real-cond-exp-ge-c[OF integrable-norm[OF
integrable], of 0] norm-ge-zero **by** *fastforce*
have **: *A \in sets F $\implies \int x \in A. |f x| \partial M = \int x \in A. \text{real-cond-exp M F } (\lambda x. \text{norm } (f x)) x \partial M$* **for** *A* **unfolding** *real-norm-def* **using** *assms integrable-abs*
real-cond-exp-intA **by** *blast*

have *norm-int: A \in sets F $\implies (\int x \in A. |f x| \partial M) = (\int^+ x \in A. |f x| \partial M)$* **for** *A*
using *assms* **by** (*intro nn-set-integral-eq-set-integral[symmetric], blast, fastforce*)
(meson subalg subalgebra-def subsetD)

have *AE x in M. real-cond-exp M F ($\lambda x. \text{norm } (f x)$) x \geq 0* **using** *int real-cond-exp-ge-c*
by *force*

hence *cond-exp-norm-int: A \in sets F $\implies (\int x \in A. \text{real-cond-exp M F } (\lambda x. \text{norm } (f x)) x \partial M) = (\int^+ x \in A. \text{real-cond-exp M F } (\lambda x. \text{norm } (f x)) x \partial M)$* **for** *A* **using**
assms **by** (*intro nn-set-integral-eq-set-integral[symmetric], blast, fastforce*) *(meson*
subalg subalgebra-def subsetD)

have *A \in sets F $\implies \int^+ x \in A. |f x| \partial M = \int^+ x \in A. \text{real-cond-exp M F } (\lambda x. \text{norm } (f x)) x \partial M$* **for** *A* **using** ** *norm-int cond-exp-norm-int* **by** (*auto simp*
add: nn-integral-set-ennreal)

moreover **have** *($\lambda x. \text{ennreal } |f x|$) \in borel-measurable M* **by** *measurable*
moreover **have** *($\lambda x. \text{ennreal } (\text{real-cond-exp M F } (\lambda x. \text{norm } (f x)) x)$) \in borel-measurable*
F **by** *measurable*

ultimately **have** *AE x in M. nn-cond-exp M F ($\lambda x. \text{ennreal } |f x|$) x = real-cond-exp*
M F ($\lambda x. \text{norm } (f x)$) x **by** (*intro nn-cond-exp-charact[THEN AE-symmetric],*
auto)

hence *AE x in M. nn-cond-exp M F ($\lambda x. \text{ennreal } |f x|$) x \leq cond-exp M F ($\lambda x. \text{norm } (f x)$) x* **using** *cond-exp-real[OF int]* **by** *force*

moreover **have** *AE x in M. |real-cond-exp M F f x| = norm (cond-exp M F f x)*
unfolding *real-norm-def* **using** *cond-exp-real[OF assms]* * **by** *force*

ultimately have $AE\ x\ in\ M. \text{ennreal}(\text{norm}(\text{cond-exp } M\ F\ f\ x)) \leq \text{cond-exp } M\ F$
 $(\lambda x. \text{norm}(f\ x))\ x$ **using** $\text{real-cond-exp-abs}[OF\ \text{assms}[THEN\ \text{borel-measurable-integrable}]]$
by fastforce
hence $AE\ x\ in\ M. \text{enn2real}(\text{ennreal}(\text{norm}(\text{cond-exp } M\ F\ f\ x))) \leq \text{enn2real}$
 $(\text{cond-exp } M\ F\ (\lambda x. \text{norm}(f\ x))\ x)$ **using** ennreal-le-iff2 **by** force
thus $?thesis$ **using** $*$ **by** fastforce
qed

lemma $\text{cond-exp-contraction-simple}$:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes $\text{simple-function } M\ f\ \text{emeasure } M\ \{y \in \text{space } M. f\ y \neq 0\} \neq \infty$
shows $AE\ x\ in\ M. \text{norm}(\text{cond-exp } M\ F\ f\ x) \leq \text{cond-exp } M\ F\ (\lambda x. \text{norm}(f\ x))\ x$
using assms
proof ($\text{induction rule: integrable-simple-function-induct}$)
case ($\text{cong } f\ g$)
hence $ae: AE\ x\ in\ M. f\ x = g\ x$ **by** blast
hence $AE\ x\ in\ M. \text{cond-exp } M\ F\ f\ x = \text{cond-exp } M\ F\ g\ x$ **using** $\text{cong has-cond-exp-simple}$
by ($\text{subst cond-exp-cong-AE}$) ($\text{auto intro!: has-cond-expD}(2)$)
hence $AE\ x\ in\ M. \text{norm}(\text{cond-exp } M\ F\ f\ x) = \text{norm}(\text{cond-exp } M\ F\ g\ x)$ **by**
 force
moreover have $AE\ x\ in\ M. \text{cond-exp } M\ F\ (\lambda x. \text{norm}(f\ x))\ x = \text{cond-exp } M\ F$
 $(\lambda x. \text{norm}(g\ x))\ x$ **using** $ae\ \text{cong has-cond-exp-simple}$ **by** ($\text{subst cond-exp-cong-AE}$)
 $(\text{auto dest: has-cond-expD})$
ultimately show $?case$ **using** $\text{cong}(6)$ **by** fastforce

next

case ($\text{indicator } A\ y$)
hence $AE\ x\ in\ M. \text{cond-exp } M\ F\ (\lambda a. \text{indicator } A\ a\ *_R\ y)\ x = \text{cond-exp } M\ F$
 $(\text{indicator } A)\ x\ *_R\ y$ **by** blast
hence $*$: $AE\ x\ in\ M. \text{norm}(\text{cond-exp } M\ F\ (\lambda a. \text{indicat-real } A\ a\ *_R\ y)\ x) \leq \text{norm } y$
 $*\ \text{cond-exp } M\ F\ (\lambda x. \text{norm}(\text{indicat-real } A\ x))\ x$ **using** $\text{cond-exp-contraction-real}[OF\ \text{integrable-real-indicator, OF indicator}]$ **by** fastforce

have $AE\ x\ in\ M. \text{norm } y * \text{cond-exp } M\ F\ (\lambda x. \text{norm}(\text{indicat-real } A\ x))\ x = \text{norm}$
 $y * \text{real-cond-exp } M\ F\ (\lambda x. \text{norm}(\text{indicat-real } A\ x))\ x$ **using** $\text{cond-exp-real}[OF\ \text{integrable-real-indicator, OF indicator}]$ **by** fastforce

moreover have $AE\ x\ in\ M. \text{cond-exp } M\ F\ (\lambda x. \text{norm } y * \text{norm}(\text{indicat-real}$
 $A\ x))\ x = \text{real-cond-exp } M\ F\ (\lambda x. \text{norm } y * \text{norm}(\text{indicat-real } A\ x))\ x$ **using**
 indicator **by** ($\text{intro cond-exp-real, auto}$)

ultimately have $AE\ x\ in\ M. \text{norm } y * \text{cond-exp } M\ F\ (\lambda x. \text{norm}(\text{indicat-real } A$
 $x))\ x = \text{cond-exp } M\ F\ (\lambda x. \text{norm } y * \text{norm}(\text{indicat-real } A\ x))\ x$ **using** $\text{real-cond-exp-cmult}[of\ \lambda x. \text{norm}(\text{indicat-real } A\ x)\ \text{norm } y]$ indicator **by** fastforce

moreover have $(\lambda x. \text{norm } y * \text{norm}(\text{indicat-real } A\ x)) = (\lambda x. \text{norm}(\text{indicat-real}$
 $A\ x\ *_R\ y))$ **by** force

ultimately show $?case$ **using** $*$ **by** force

next

case ($\text{add } u\ v$)

have $AE\ x\ in\ M. \text{norm}(\text{cond-exp } M\ F\ (\lambda a. u\ a + v\ a)\ x) = \text{norm}(\text{cond-exp } M$
 $F\ u\ x + \text{cond-exp } M\ F\ v\ x)$ **using** $\text{has-cond-exp-charact}(2)[OF\ \text{has-cond-exp-add, OF has-cond-exp-simple}(1,1), OF\ \text{add}(1,2,3,4)]$ **by** fastforce

moreover have $AE\ x\ in\ M.\ norm\ (cond\text{-}exp\ M\ F\ u\ x + cond\text{-}exp\ M\ F\ v\ x) \leq$
 $norm\ (cond\text{-}exp\ M\ F\ u\ x) + norm\ (cond\text{-}exp\ M\ F\ v\ x)$ **using** *norm-triangle-ineq*
by *blast*
moreover have $AE\ x\ in\ M.\ norm\ (cond\text{-}exp\ M\ F\ u\ x) + norm\ (cond\text{-}exp\ M\ F\ v\ x) \leq$
 $cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (u\ x))\ x + cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (v\ x))\ x$ **using**
 $add(6,7)$ **by** *fastforce*
moreover have $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (u\ x))\ x + cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (v\ x))\ x =$
 $cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (u\ x) + norm\ (v\ x))\ x$ **using** *integrable-simple-function[OF add(1,2)] integrable-simple-function[OF add(3,4)]* **by**
 $(intro\ has\text{-}cond\text{-}exp\text{-}charact(2)[OF\ has\text{-}cond\text{-}exp\text{-}add[OF\ has\text{-}cond\text{-}exp\text{-}charact(1,1)],$
 $THEN\ AE\text{-}symmetric],\ auto\ intro:\ has\text{-}cond\text{-}exp\text{-}real)$
moreover have $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (u\ x) + norm\ (v\ x))\ x =$
 $cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (u\ x + v\ x))\ x$ **using** $add(5)$ *integrable-simple-function[OF add(1,2)] integrable-simple-function[OF add(3,4)]* **by** $(intro\ cond\text{-}exp\text{-}cong,\ auto)$
ultimately show *?case* **by** *force*
qed

lemma *has-cond-exp-simple-lim*:

fixes $f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology,\ banach\}$
assumes *integrable[measurable]: integrable M f*
and $\bigwedge i.\ simple\text{-}function\ M\ (s\ i)$
and $\bigwedge i.\ emeasure\ M\ \{y \in space\ M.\ s\ i\ y \neq 0\} \neq \infty$
and $\bigwedge x.\ x \in space\ M \implies (\lambda i.\ s\ i\ x) \longrightarrow f\ x$
and $\bigwedge x\ i.\ x \in space\ M \implies norm\ (s\ i\ x) \leq 2 * norm\ (f\ x)$
obtains r
where *strict-mono r has-cond-exp M F f* $(\lambda x.\ lim\ (\lambda i.\ cond\text{-}exp\ M\ F\ (s\ (r\ i))\ x))$
 $AE\ x\ in\ M.\ convergent\ (\lambda i.\ cond\text{-}exp\ M\ F\ (s\ (r\ i))\ x)$
proof –
have *[measurable]: (s i) ∈ borel-measurable M* **for** i **using** *assms(2)* **by** *(simp add: borel-measurable-simple-function)*
have *integrable-s: integrable M (λx. s i x)* **for** i **using** *assms integrable-simple-function*
by *blast*
have *integrable-4f: integrable M (λx. 4 * norm (f x))* **using** *assms(1)* **by** *simp*
have *integrable-2f: integrable M (λx. 2 * norm (f x))* **using** *assms(1)* **by** *simp*
have *integrable-2-cond-exp-norm-f: integrable M (λx. 2 * cond-exp M F (λx. norm (f x)) x)* **by** *fast*

have $emeasure\ M\ \{y \in space\ M.\ s\ i\ y - s\ j\ y \neq 0\} \leq emeasure\ M\ \{y \in space\ M.\ s\ i\ y \neq 0\} + emeasure\ M\ \{y \in space\ M.\ s\ j\ y \neq 0\}$ **for** $i\ j$ **using**
 $simple\text{-}functionD(2)[OF\ assms(2)]$ **by** $(intro\ order\text{-}trans[OF\ emeasure\text{-}mono\ emeasure\text{-}subadditive],\ auto)$

hence *fin-sup: emeasure M {y ∈ space M. s i y - s j y ≠ 0} ≠ ∞* **for** $i\ j$ **using** *assms(3)* **by** $(metis\ (mono\text{-}tags)\ ennreal\text{-}add\text{-}eq\text{-}top\ linorder\text{-}not\text{-}less\ top.\ not\text{-}eq\text{-}extremum\ infinity\text{-}ennreal\text{-}def)$

have $emeasure\ M\ \{y \in space\ M.\ norm\ (s\ i\ y - s\ j\ y) \neq 0\} \leq emeasure\ M\ \{y \in space\ M.\ s\ i\ y \neq 0\} + emeasure\ M\ \{y \in space\ M.\ s\ j\ y \neq 0\}$ **for** $i\ j$ **using**
 $simple\text{-}functionD(2)[OF\ assms(2)]$ **by** $(intro\ order\text{-}trans[OF\ emeasure\text{-}mono]$

emeasure-subadditive], *auto*)

hence *fin-sup-norm*: *emeasure* $M \{y \in \text{space } M. \text{norm } (s \ i \ y - s \ j \ y) \neq 0\} \neq \infty$
for $i \ j$ **using** *assms*(3) **by** (*metis* (*mono-tags*) *ennreal-add-eq-top linorder-not-less top.not-eq-extremum infinity-ennreal-def*)

have *Cauchy*: *Cauchy* $(\lambda n. s \ n \ x)$ **if** $x \in \text{space } M$ **for** x **using** *assms*(4) *LIM-SEQ-imp-Cauchy* **that by** *blast*

hence *bounded-range-s*: *bounded* $(\text{range } (\lambda n. s \ n \ x))$ **if** $x \in \text{space } M$ **for** x **using** *that cauchy-imp-bounded by fast*

have *AE* x **in** $M. (\lambda n. \text{diameter } \{s \ i \ x \mid i. n \leq i\}) \longrightarrow 0$ **using** *Cauchy cauchy-iff-diameter-tends-to-zero-and-bounded by fast*

moreover **have** $(\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\}) \in \text{borel-measurable } M$ **for** n **using** *bounded-range-s borel-measurable-diameter by measurable*

moreover **have** *AE* x **in** $M. \text{norm } (\text{diameter } \{s \ i \ x \mid i. n \leq i\}) \leq 4 * \text{norm } (f \ x)$ **for** n

proof –

{

fix x **assume** $x: x \in \text{space } M$

have $\text{diameter } \{s \ i \ x \mid i. n \leq i\} \leq 2 * \text{norm } (f \ x) + 2 * \text{norm } (f \ x)$

by (*intro diameter-le, blast, subst dist-norm[symmetric], intro dist-triangle3[THEN order-trans, of 0], intro add-mono*) (*auto intro: assms*(5)[*OF* x])

hence $\text{norm } (\text{diameter } \{s \ i \ x \mid i. n \leq i\}) \leq 4 * \text{norm } (f \ x)$ **using** *diameter-ge-0[OF bounded-subset[OF bounded-range-s], OF x, of {s i x | i. n ≤ i}] by force*

}

thus *?thesis by fast*

qed

ultimately **have** *diameter-tendsto-zero*: $(\lambda n. \text{LINT } x | M. \text{diameter } \{s \ i \ x \mid i. n \leq i\}) \longrightarrow 0$ **by** (*intro integral-dominated-convergence[OF borel-measurable-const[of 0] - integrable-4f, simplified]*) (*fast+*)

have *diameter-integrable*: *integrable* $M (\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\})$ **for** n **using** *assms*(1,5)

by (*intro integrable-bound-diameter[OF bounded-range-s integrable-2f], auto*)

have *dist-integrable*: *integrable* $M (\lambda x. \text{dist } (s \ i \ x) (s \ j \ x))$ **for** $i \ j$ **using** *assms*(5) *dist-triangle3[of s i - 0, THEN order-trans, OF add-mono, of - 2 * norm (f -)]*

by (*intro Bochner-Integration.integrable-bound[OF integrable-4f]*) *fastforce+*

have $\exists N. \forall i \geq N. \forall j \geq N. \text{LINT } x | M. \text{norm } (\text{cond-exp } M \ F \ (s \ i) \ x - \text{cond-exp } M \ F \ (s \ j) \ x) < e$ **if** $e\text{-pos}: e > 0$ **for** e

proof –

obtain N **where** $*$: $\text{LINT } x | M. \text{diameter } \{s \ i \ x \mid i. n \leq i\} < e$ **if** $n \geq N$ **for** n **using** *that order-tendsto-iff[THEN iffD1, OF diameter-tendsto-zero, unfolded eventually-sequentially] e-pos by presburger*

{

fix $i \ j \ x$ **assume** *asm*: $i \geq N \ j \geq N \ x \in \text{space } M$

have *case-prod dist* ‘ $(\{s \ i \ x \mid i. N \leq i\} \times \{s \ i \ x \mid i. N \leq i\}) = \text{case-prod } (\lambda i$

$j. \text{ dist } (s \ i \ x) \ (s \ j \ x)) \text{ ' } (\{N..\} \times \{N..\}) \text{ by fast}$
hence $\text{diameter } \{s \ i \ x \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}). \text{ dist } (s \ i \ x) \ (s \ j \ x))$ **unfolding** diameter-def **by** auto
moreover **have** $(\text{SUP } (i, j) \in \{N..\} \times \{N..\}). \text{ dist } (s \ i \ x) \ (s \ j \ x)) \geq \text{dist } (s \ i \ x) \ (s \ j \ x)$ **using** $\text{asm bounded-imp-bdd-above[OF bounded-imp-dist-bounded, OF bounded-range-s]}$ **by** $(\text{intro cSup-upper, auto})$
ultimately **have** $\text{diameter } \{s \ i \ x \mid i. N \leq i\} \geq \text{dist } (s \ i \ x) \ (s \ j \ x)$ **by** presburger
}
hence $\text{LINT } x \mid M. \text{ dist } (s \ i \ x) \ (s \ j \ x) < e \text{ if } i \geq N \ j \geq N \text{ for } i \ j$ **using** $\text{that * by (intro integral-mono[OF dist-integrable diameter-integrable, THEN order.strict-trans1], blast+)}$
moreover **have** $\text{LINT } x \mid M. \text{ norm } (\text{cond-exp } M \ F \ (s \ i) \ x - \text{cond-exp } M \ F \ (s \ j) \ x) \leq \text{LINT } x \mid M. \text{ dist } (s \ i \ x) \ (s \ j \ x)$ **for** $i \ j$
proof $-$
have $\text{LINT } x \mid M. \text{ norm } (\text{cond-exp } M \ F \ (s \ i) \ x - \text{cond-exp } M \ F \ (s \ j) \ x) = \text{LINT } x \mid M. \text{ norm } (\text{cond-exp } M \ F \ (s \ i) \ x + -1 * \text{cond-exp } M \ F \ (s \ j) \ x)$ **unfolding** dist-norm **by** simp
also **have** $\dots = \text{LINT } x \mid M. \text{ norm } (\text{cond-exp } M \ F \ (\lambda x. s \ i \ x - s \ j \ x) \ x)$ **using** $\text{has-cond-exp-charact(2)[OF has-cond-exp-add[OF - has-cond-exp-scaleR-right, OF has-cond-exp-charact(1,1), OF has-cond-exp-simple(1,1)[OF assms(2,3)]]}$, $\text{THEN AE-symmetric, of } i - 1 \ j]$ **by** $(\text{intro integral-cong-AE}) \text{ force+}$
also **have** $\dots \leq \text{LINT } x \mid M. \text{ cond-exp } M \ F \ (\lambda x. \text{norm } (s \ i \ x - s \ j \ x)) \ x$ **using** $\text{cond-exp-contraction-simple[OF - fin-sup, of } i \ j]$ $\text{integrable-cond-exp assms(2)}$ **by** $(\text{intro integral-mono-AE, fast+})$
also **have** $\dots = \text{LINT } x \mid M. \text{ norm } (s \ i \ x - s \ j \ x)$ **unfolding** $\text{set-integral-space(1)[OF integrable-cond-exp, symmetric]}$ $\text{set-integral-space[OF dist-integrable[unfolding dist-norm], symmetric]}$ **by** $(\text{intro has-cond-expD(1)[OF has-cond-exp-simple[OF - fin-sup-norm], symmetric]})$ $(\text{metis assms(2) simple-function-compose1 simple-function-diff, metis sets.top subalg subalgebra-def})$
finally **show** $?thesis$ **unfolding** dist-norm .
qed
ultimately **show** $?thesis$ **using** $\text{order.strict-trans1}$ **by** meson
qed
then **obtain** r **where** $\text{strict-mono-r: strict-mono } r$ **and** $\text{AE-Cauchy: AE } x \text{ in } M. \text{ Cauchy } (\lambda i. \text{cond-exp } M \ F \ (s \ (r \ i)) \ x)$ **by** $(\text{rule cauchy-L1-AE-cauchy-subseq[OF integrable-cond-exp], auto})$
hence $\text{ae-lim-cond-exp: AE } x \text{ in } M. (\lambda n. \text{cond-exp } M \ F \ (s \ (r \ n)) \ x) \longrightarrow \lim (\lambda n. \text{cond-exp } M \ F \ (s \ (r \ n)) \ x)$ **using** $\text{Cauchy-convergent-iff convergent-LIMSEQ-iff}$ **by** fastforce

have $\text{cond-exp-bounded: AE } x \text{ in } M. \text{ norm } (\text{cond-exp } M \ F \ (s \ (r \ n)) \ x) \leq \text{cond-exp } M \ F \ (\lambda x. 2 * \text{norm } (f \ x)) \ x$ **for** n
proof $-$
have $\text{AE } x \text{ in } M. \text{ norm } (\text{cond-exp } M \ F \ (s \ (r \ n)) \ x) \leq \text{cond-exp } M \ F \ (\lambda x. \text{norm } (s \ (r \ n) \ x)) \ x$ **by** $(\text{rule cond-exp-contraction-simple[OF assms(2,3)])}$
moreover **have** $\text{AE } x \text{ in } M. \text{ real-cond-exp } M \ F \ (\lambda x. \text{norm } (s \ (r \ n) \ x)) \ x \leq \text{real-cond-exp } M \ F \ (\lambda x. 2 * \text{norm } (f \ x)) \ x$ **using** $\text{integrable-s integrable-2f assms(5)}$ **by** $(\text{intro real-cond-exp-mono, auto})$

ultimately show *?thesis* using *cond-exp-real*[*OF integrable-norm*, *OF integrable-s*, of *r n*] *cond-exp-real*[*OF integrable-2f*] by force

qed

have *lim-integrable*: *integrable* *M* ($\lambda x. \lim (\lambda i. \text{cond-exp } M \ F \ (s \ (r \ i)) \ x)$)

by (intro *integrable-dominated-convergence*[*OF - borel-measurable-cond-exp'* *integrable-cond-exp ae-lim-cond-exp cond-exp-bounded*], *simp*)

{

fix *A* assume *A-in-sets-F*: *A* \in *sets F*

have *AE* *x* in *M*. *norm* (*indicator* *A* *x* \ast_R *cond-exp* *M* *F* (*s* (*r* *n*)) *x*) \leq *cond-exp* *M* *F* ($\lambda x. 2 \ast \text{norm} \ (f \ x)$) *x* for *n*

proof -

have *AE* *x* in *M*. *norm* (*indicator* *A* *x* \ast_R *cond-exp* *M* *F* (*s* (*r* *n*)) *x*) \leq *norm* (*cond-exp* *M* *F* (*s* (*r* *n*)) *x*) unfolding *indicator-def* by *simp*

thus *?thesis* using *cond-exp-bounded*[of *n*] by force

qed

hence *lim-cond-exp-int*: ($\lambda n. \text{LINT } x:A|M. \text{cond-exp } M \ F \ (s \ (r \ n)) \ x$) \longrightarrow *LINT* *x:A*|*M*. $\lim (\lambda n. \text{cond-exp } M \ F \ (s \ (r \ n)) \ x)$

using *ae-lim-cond-exp measurable-from-subalg*[*OF subalg borel-measurable-indicator*, *OF A-in-sets-F*] *cond-exp-bounded*

unfolding *set-lebesgue-integral-def*

by (intro *integral-dominated-convergence*[*OF borel-measurable-scaleR borel-measurable-scaleR integrable-cond-exp*]) (*fastforce simp add: tendsto-scaleR*) +

have *AE* *x* in *M*. *norm* (*indicator* *A* *x* \ast_R *s* (*r* *n*) *x*) $\leq 2 \ast \text{norm} \ (f \ x)$ for *n*

proof -

have *AE* *x* in *M*. *norm* (*indicator* *A* *x* \ast_R *s* (*r* *n*) *x*) \leq *norm* (*s* (*r* *n*) *x*)

unfolding *indicator-def* by *simp*

thus *?thesis* using *assms*(5)[of - *r n*] by *fastforce*

qed

hence *lim-s-int*: ($\lambda n. \text{LINT } x:A|M. \text{s} \ (r \ n) \ x$) \longrightarrow *LINT* *x:A*|*M*. *f* *x*

using *measurable-from-subalg*[*OF subalg borel-measurable-indicator*, *OF A-in-sets-F*] *LIMSEQ-subseq-LIMSEQ*[*OF assms*(4) *strict-mono-r*] *assms*(5)

unfolding *set-lebesgue-integral-def comp-def*

by (intro *integral-dominated-convergence*[*OF borel-measurable-scaleR borel-measurable-scaleR integrable-2f*]) (*fastforce simp add: tendsto-scaleR*) +

have *LINT* *x:A*|*M*. $\lim (\lambda n. \text{cond-exp } M \ F \ (s \ (r \ n)) \ x) = \lim (\lambda n. \text{LINT } x:A|M. \text{cond-exp } M \ F \ (s \ (r \ n)) \ x)$ using *limI*[*OF lim-cond-exp-int*] by *argo*

also have ... = $\lim (\lambda n. \text{LINT } x:A|M. \text{s} \ (r \ n) \ x)$ using *has-cond-expD*(1)[*OF has-cond-exp-simple*[*OF assms*(2,3)] *A-in-sets-F*, *symmetric*] by *presburger*

also have ... = *LINT* *x:A*|*M*. *f* *x* using *limI*[*OF lim-s-int*] by *argo*

finally have *LINT* *x:A*|*M*. $\lim (\lambda n. \text{cond-exp } M \ F \ (s \ (r \ n)) \ x) = \text{LINT } x:A|M. \text{f } x$.

}

hence *has-cond-exp* *M* *F* *f* ($\lambda x. \lim (\lambda i. \text{cond-exp } M \ F \ (s \ (r \ i)) \ x)$) using *assms*(1) *lim-integrable* by (intro *has-cond-expI'*, *auto*)

thus *thesis* using *AE-Cauchy Cauchy-convergent strict-mono-r* by (*auto intro!*: *that*)

qed

corollary *has-cond-expI*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f$
shows *has-cond-exp* $M F f$ (*cond-exp* $M F f$)
proof –
obtain s **where** $s\text{-is}$: $\bigwedge i. \text{simple-function } M (s i) \bigwedge i. \text{emeasure } M \{y \in \text{space } M. s i y \neq 0\} \neq \infty \bigwedge x. x \in \text{space } M \implies (\lambda i. s i x) \longrightarrow f x \bigwedge x i. x \in \text{space } M \implies \text{norm } (s i x) \leq 2 * \text{norm } (f x)$ **using** *integrable-implies-simple-function-sequence*[*OF assms*] **by** *blast*
show *?thesis* **using** *has-cond-exp-simple-lim*[*OF assms s-is*] *has-cond-exp-charact*(1)
by *metis*
qed

lemma *cond-exp-nested-subalg*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f$ *subalgebra* $M G$ *subalgebra* $G F$
shows $AE \xi$ *in* $M. \text{cond-exp } M F f \xi = \text{cond-exp } M F (\text{cond-exp } M G f) \xi$
using *has-cond-expI* *assms sigma-finite-subalgebra-def* **by** (*auto intro!*: *has-cond-exp-nested-subalg*[*THEN has-cond-exp-charact*(2), *THEN AE-symmetric*] *sigma-finite-subalgebra.has-cond-expI*[*OF sigma-finite-subalgebra.intro*[*OF assms*(2)]]] *nested-subalg-is-sigma-finite*)

lemma *cond-exp-set-integral*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f$ $A \in \text{sets } F$
shows $(\int x \in A. f x \partial M) = (\int x \in A. \text{cond-exp } M F f x \partial M)$
using *has-cond-expD*(1)[*OF has-cond-expI*, *OF assms*] **by** *argo*

lemma *cond-exp-add*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f$ *integrable* $M g$
shows $AE x$ *in* $M. \text{cond-exp } M F (\lambda x. f x + g x) x = \text{cond-exp } M F f x + \text{cond-exp } M F g x$
using *has-cond-exp-add*[*OF has-cond-expI*(1,1), *OF assms*, *THEN has-cond-exp-charact*(2)]
.

lemma *cond-exp-diff*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f$ *integrable* $M g$
shows $AE x$ *in* $M. \text{cond-exp } M F (\lambda x. f x - g x) x = \text{cond-exp } M F f x - \text{cond-exp } M F g x$
using *has-cond-exp-add*[*OF - has-cond-exp-scaleR-right*, *OF has-cond-expI*(1,1), *OF assms*, *THEN has-cond-exp-charact*(2), *of -1*] **by** *simp*

lemma *cond-exp-diff'*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$

assumes *integrable M f integrable M g*
shows $AE\ x\ in\ M.\ cond_exp\ M\ F\ (f - g)\ x = cond_exp\ M\ F\ f\ x - cond_exp\ M\ F\ g\ x$
unfolding *fun-diff-def* **using** *assms* **by** (*rule cond-exp-diff*)

lemma *cond-exp-scaleR-left*:

fixes $f :: 'a \Rightarrow real$
assumes *integrable M f*
shows $AE\ x\ in\ M.\ cond_exp\ M\ F\ (\lambda x.\ f\ x *_{\mathbb{R}} c)\ x = cond_exp\ M\ F\ f\ x *_{\mathbb{R}} c$
using *cond-exp-set-integral[OF assms]* *subalg assms* **unfolding** *subalgebra-def*
by (*intro cond-exp-charact,*
subst set-integral-scaleR-left, blast, intro assms,
subst set-integral-scaleR-left, blast, intro integrable-cond-exp)
auto

lemma *cond-exp-contraction*:

fixes $f :: 'a \Rightarrow 'b :: \{second-countable-topology, banach\}$
assumes *integrable M f*
shows $AE\ x\ in\ M.\ norm\ (cond_exp\ M\ F\ f\ x) \leq cond_exp\ M\ F\ (\lambda x.\ norm\ (f\ x))\ x$
proof –

obtain s **where** $s: \bigwedge i.\ simple_function\ M\ (s\ i) \bigwedge i.\ emeasure\ M\ \{y \in space\ M.\ s\ i\ y \neq 0\} \neq \infty \bigwedge x.\ x \in space\ M \implies (\lambda i.\ s\ i\ x) \longrightarrow f\ x \bigwedge i.\ x \in space\ M \implies norm\ (s\ i\ x) \leq 2 * norm\ (f\ x)$
by (*blast intro: integrable-implies-simple-function-sequence[OF assms]*)

obtain r **where** $r: strict_mono\ r$ **and** *has-cond-exp M F f* $(\lambda x.\ lim\ (\lambda i.\ cond_exp\ M\ F\ (s\ (r\ i))\ x))\ AE\ x\ in\ M.\ (\lambda i.\ cond_exp\ M\ F\ (s\ (r\ i))\ x) \longrightarrow lim\ (\lambda i.\ cond_exp\ M\ F\ (s\ (r\ i))\ x)$

using *has-cond-exp-simple-lim[OF assms s]* **unfolding** *convergent-LIMSEQ-iff*
by *blast*

hence *r-tendsto*: $AE\ x\ in\ M.\ (\lambda i.\ cond_exp\ M\ F\ (s\ (r\ i))\ x) \longrightarrow cond_exp\ M\ F\ f\ x$ **using** *has-cond-exp-charact(2)* **by** *force*

have *norm-s-r*: $\bigwedge i.\ simple_function\ M\ (\lambda x.\ norm\ (s\ (r\ i)\ x)) \bigwedge i.\ emeasure\ M\ \{y \in space\ M.\ norm\ (s\ (r\ i)\ y) \neq 0\} \neq \infty \bigwedge x.\ x \in space\ M \implies (\lambda i.\ norm\ (s\ (r\ i)\ x)) \longrightarrow norm\ (f\ x) \bigwedge i.\ x \in space\ M \implies norm\ (norm\ (s\ (r\ i)\ x)) \leq 2 * norm\ (norm\ (f\ x))$

using s **by** (*auto intro: LIMSEQ-subseq-LIMSEQ[OF tendsto-norm r, unfolded comp-def]* *simple-function-compose1*)

obtain r' **where** $r': strict_mono\ r'$ **and** *has-cond-exp M F* $(\lambda x.\ norm\ (f\ x))\ (\lambda x.\ lim\ (\lambda i.\ cond_exp\ M\ F\ (\lambda x.\ norm\ (s\ (r\ (r'\ i))\ x))\ x))\ AE\ x\ in\ M.\ (\lambda i.\ cond_exp\ M\ F\ (\lambda x.\ norm\ (s\ (r\ (r'\ i))\ x))\ x) \longrightarrow lim\ (\lambda i.\ cond_exp\ M\ F\ (\lambda x.\ norm\ (s\ (r\ (r'\ i))\ x))\ x)$ **using** *has-cond-exp-simple-lim[OF integrable-norm norm-s-r, OF assms]* **unfolding** *convergent-LIMSEQ-iff* **by** *blast*

hence *r'-tendsto*: $AE\ x\ in\ M.\ (\lambda i.\ cond_exp\ M\ F\ (\lambda x.\ norm\ (s\ (r\ (r'\ i))\ x))\ x) \longrightarrow cond_exp\ M\ F\ (\lambda x.\ norm\ (f\ x))\ x$ **using** *has-cond-exp-charact(2)* **by** *force*

have $AE\ x\ in\ M. \forall i. norm\ (cond-exp\ M\ F\ (s\ (r\ (r'\ i)))\ x) \leq cond-exp\ M\ F\ (\lambda x. norm\ (s\ (r\ (r'\ i)))\ x)\ x$ **using** s **by** (*auto intro: cond-exp-contraction-simple simp add: AE-all-countable*)
moreover have $AE\ x\ in\ M. (\lambda i. norm\ (cond-exp\ M\ F\ (s\ (r\ (r'\ i)))\ x)) \longrightarrow norm\ (cond-exp\ M\ F\ f\ x)$ **using** $r-tendsto\ LIMSEQ-subseq-LIMSEQ[OF\ tendsto-norm\ r',\ unfolded\ comp-def]$ **by** *fast*
ultimately show *?thesis* **using** $LIMSEQ-le\ r'-tendsto$ **by** *fast*
qed

lemma *cond-exp-measurable-mult:*

fixes $f\ g :: 'a \Rightarrow real$
assumes [*measurable*]: $integrable\ M\ (\lambda x. f\ x * g\ x)\ integrable\ M\ g\ f \in borel-measurable\ F$

shows $integrable\ M\ (\lambda x. f\ x * cond-exp\ M\ F\ g\ x)$

$AE\ x\ in\ M. cond-exp\ M\ F\ (\lambda x. f\ x * g\ x)\ x = f\ x * cond-exp\ M\ F\ g\ x$

proof –

show *integrable: integrable* $M\ (\lambda x. f\ x * cond-exp\ M\ F\ g\ x)$ **using** *cond-exp-real*[*OF* *assms*(2)] **by** (*intro integrable-cong-AE-imp*[*OF* *real-cond-exp-intg*(1), *OF* *assms*(1,3) *assms*(2)[*THEN borel-measurable-integrable*]] *measurable-from-subalg*[*OF* *subalg*])
auto

interpret *sigma-finite-measure restr-to-subalg* $M\ F$ **by** (*rule sigma-fin-subalg*)

{

fix A **assume** *asm*: $A \in sets\ F$

hence *asm'*: $A \in sets\ M$ **using** *subalg* **by** (*fastforce simp add: subalgebra-def*)

have *set-lebesgue-integral* $M\ A\ (cond-exp\ M\ F\ (\lambda x. f\ x * g\ x)) = set-lebesgue-integral\ M\ A\ (\lambda x. f\ x * g\ x)$ **by** (*simp add: cond-exp-set-integral*[*OF* *assms*(1) *asm*])

also have $\dots = set-lebesgue-integral\ M\ A\ (\lambda x. f\ x * real-cond-exp\ M\ F\ g\ x)$ **using** *borel-measurable-times*[*OF* *borel-measurable-indicator*[*OF* *asm*] *assms*(3)]
borel-measurable-integrable[*OF* *assms*(2)] *integrable-mult-indicator*[*OF* *asm'* *assms*(1)]

by (*fastforce simp add: set-lebesgue-integral-def mult.assoc*[*symmetric*] *intro: real-cond-exp-intg*(2)[*symmetric*])

also have $\dots = set-lebesgue-integral\ M\ A\ (\lambda x. f\ x * cond-exp\ M\ F\ g\ x)$ **using**
cond-exp-real[*OF* *assms*(2)] *asm'* *borel-measurable-cond-exp'* *borel-measurable-cond-exp2*
measurable-from-subalg[*OF* *subalg* *assms*(3)] **by** (*auto simp add: set-lebesgue-integral-def intro: integral-cong-AE*)

finally have *set-lebesgue-integral* $M\ A\ (cond-exp\ M\ F\ (\lambda x. f\ x * g\ x)) = \int_{x \in A} f\ x * g\ x\ dM$.

}

hence $AE\ x\ in\ restr-to-subalg\ M\ F. cond-exp\ M\ F\ (\lambda x. f\ x * g\ x)\ x = f\ x * cond-exp\ M\ F\ g\ x$ **by** (*intro density-unique-banach integrable-cond-exp integrable-integrable-in-subalg subalg, measurable, simp add: set-lebesgue-integral-def integral-subalgebra2*[*OF* *subalg*] *sets-restr-to-subalg*[*OF* *subalg*])

thus $AE\ x\ in\ M. cond-exp\ M\ F\ (\lambda x. f\ x * g\ x)\ x = f\ x * cond-exp\ M\ F\ g\ x$ **by**
(rule AE-restr-to-subalg[*OF* *subalg*])

qed

lemma *cond-exp-measurable-scaleR:*

fixes $f :: 'a \Rightarrow real$ **and** $g :: 'a \Rightarrow 'b :: \{second-countable-topology, banach\}$

assumes $[measurable]: \text{integrable } M (\lambda x. f x *_{\mathbb{R}} g x) \text{ integrable } M g f \in \text{borel-measurable } F$

shows $\text{integrable } M (\lambda x. f x *_{\mathbb{R}} \text{cond-exp } M F g x)$

$AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x *_{\mathbb{R}} g x) x = f x *_{\mathbb{R}} \text{cond-exp } M F g x$

proof –

let $?F = \text{restr-to-subalg } M F$

have $\text{subalg}' : \text{subalgebra } M (\text{restr-to-subalg } M F) \text{ by } (\text{metis sets-eq-imp-space-eq sets-restr-to-subalg subalg subalgebra-def})$

{

fix z **assume** $\text{asm}[measurable]: \text{integrable } M (\lambda x. z x *_{\mathbb{R}} g x) z \in \text{borel-measurable } ?F$

hence $\text{asm}'[measurable]: z \in \text{borel-measurable } F$ **using** $\text{measurable-in-subalg}'$ **subalg by blast**

have $\text{integrable } M (\lambda x. z x *_{\mathbb{R}} \text{cond-exp } M F g x) \text{ LINT } x | M. z x *_{\mathbb{R}} g x = \text{LINT } x | M. z x *_{\mathbb{R}} \text{cond-exp } M F g x$

proof –

obtain s **where** $s\text{-is}: \bigwedge i. \text{simple-function } ?F (s i) \bigwedge x. x \in \text{space } ?F \implies (\lambda i. s i x) \longrightarrow z x \bigwedge i x. x \in \text{space } ?F \implies \text{norm } (s i x) \leq 2 * \text{norm } (z x)$ **using** $\text{borel-measurable-implies-sequence-metric}[OF \text{asm}(2), \text{of } 0]$ **by force**

have $s\text{-scaleR-g-tendsto}: AE x \text{ in } M. (\lambda i. s i x *_{\mathbb{R}} g x) \longrightarrow z x *_{\mathbb{R}} g x$ **using** $s\text{-is}(2)$ **by** $(\text{simp add: space-restr-to-subalg tendsto-scaleR})$

have $s\text{-scaleR-cond-exp-g-tendsto}: AE x \text{ in } ?F. (\lambda i. s i x *_{\mathbb{R}} \text{cond-exp } M F g x) \longrightarrow z x *_{\mathbb{R}} \text{cond-exp } M F g x$ **using** $s\text{-is}(2)$ **by** $(\text{simp add: tendsto-scaleR})$

have $s\text{-scaleR-g-meas}: (\lambda x. s i x *_{\mathbb{R}} g x) \in \text{borel-measurable } M$ **for** i **using** $s\text{-is}(1)[\text{THEN borel-measurable-simple-function, THEN subalg}'[\text{THEN measurable-from-subalg}]]$ **by simp**

have $s\text{-scaleR-cond-exp-g-meas}: (\lambda x. s i x *_{\mathbb{R}} \text{cond-exp } M F g x) \in \text{borel-measurable } ?F$ **for** i **using** $s\text{-is}(1)[\text{THEN borel-measurable-simple-function}]$ $\text{measurable-in-subalg}[OF \text{subalg borel-measurable-cond-exp}]$ **by** $(\text{fastforce intro: borel-measurable-scaleR})$

have $s\text{-scaleR-g-AE-bdd}: AE x \text{ in } M. \text{norm } (s i x *_{\mathbb{R}} g x) \leq 2 * \text{norm } (z x *_{\mathbb{R}} g x)$ **for** i **using** $s\text{-is}(3)$ **by** $(\text{fastforce simp add: space-restr-to-subalg mult.assoc[symmetric] mult-right-mono})$

{

fix i

have $\text{asm}: \text{integrable } M (\lambda x. \text{norm } (z x) * \text{norm } (g x))$ **using** $\text{asm}(1)[\text{THEN integrable-norm}]$ **by simp**

have $AE x \text{ in } ?F. \text{norm } (s i x *_{\mathbb{R}} \text{cond-exp } M F g x) \leq 2 * \text{norm } (z x) * \text{norm } (\text{cond-exp } M F g x)$ **using** $s\text{-is}(3)$ **by** $(\text{fastforce simp add: mult-mono})$

moreover have $AE x \text{ in } ?F. \text{norm } (z x) * \text{cond-exp } M F (\lambda x. \text{norm } (g x)) x = \text{cond-exp } M F (\lambda x. \text{norm } (z x) * \text{norm } (g x)) x$ **by** $(\text{rule cond-exp-measurable-mult}(2)[\text{THEN AE-symmetric, OF asm integrable-norm, OF assms}(2), \text{THEN AE-restr-to-subalg2}[OF \text{subalg}], \text{auto})$

ultimately have $AE x \text{ in } ?F. \text{norm } (s i x *_{\mathbb{R}} \text{cond-exp } M F g x) \leq 2 * \text{cond-exp } M F (\lambda x. \text{norm } (z x *_{\mathbb{R}} g x)) x$ **using** $\text{cond-exp-contraction}[OF \text{assms}(2),$

THEN $AE\text{-restr-to-subalg2}[OF\ subalg]$ order-trans[$OF - mult\text{-}mono$] **by** fastforce
 }
 note $s\text{-scaleR-cond-exp-g-AE-bdd} = this$

{
 fix i
 have $s\text{-meas-}M[measurable]: s\ i \in \text{borel-measurable } M$ **by** (meson borel-measurable-simple-function
 measurable-from-subalg $s\text{-is}(1)$ subalg')
 have $s\text{-meas-}F[measurable]: s\ i \in \text{borel-measurable } F$ **by** (meson borel-measurable-simple-function
 measurable-in-subalg' $s\text{-is}(1)$ subalg)

 have $s\text{-scaleR-eq}: s\ i\ x *_R h\ x = (\sum y \in s\ i\ \text{'space } M. (\text{indicator } (s\ i\ -\ \{y\} \cap \text{space } M)\ x *_R y) *_R h\ x)$ **if** $x \in \text{space } M$ **for** x **and** $h :: 'a \Rightarrow 'b$
using simple-function-indicator-representation[$OF\ s\text{-is}(1)$, of $x\ i$] that unfolding
 space-restr-to-subalg scaleR-left.sum[of - - $h\ x$, symmetric] **by** presburger

 have $LINT\ x|M. s\ i\ x *_R g\ x = LINT\ x|M. (\sum y \in s\ i\ \text{'space } M. \text{indicator } (s\ i\ -\ \{y\} \cap \text{space } M)\ x *_R y *_R g\ x)$ **using** $s\text{-scaleR-eq}$ **by** (intro Bochner-Integration.integral-cong) auto
 also have $\dots = (\sum y \in s\ i\ \text{'space } M. LINT\ x|M. \text{indicator } (s\ i\ -\ \{y\} \cap \text{space } M)\ x *_R y *_R g\ x)$ **by** (intro Bochner-Integration.integral-sum integrable-mult-indicator[$OF - integrable\text{-}scaleR\text{-}right$] assms(2)) simp
 also have $\dots = (\sum y \in s\ i\ \text{'space } M. y *_R \text{set-lebesgue-integral } M\ (s\ i\ -\ \{y\} \cap \text{space } M)\ g)$ **by** (simp only: set-lebesgue-integral-def[symmetric]) simp
 also have $\dots = (\sum y \in s\ i\ \text{'space } M. y *_R \text{set-lebesgue-integral } M\ (s\ i\ -\ \{y\} \cap \text{space } M)\ (cond\text{-exp } M\ F\ g))$ **using** assms(2) subalg borel-measurable-vimage[$OF\ s\text{-meas-}F$] **by** (subst cond-exp-set-integral, auto simp add: subalgebra-def)
 also have $\dots = (\sum y \in s\ i\ \text{'space } M. LINT\ x|M. \text{indicator } (s\ i\ -\ \{y\} \cap \text{space } M)\ x *_R y *_R cond\text{-exp } M\ F\ g\ x)$ **by** (simp only: set-lebesgue-integral-def[symmetric]) simp
 also have $\dots = LINT\ x|M. (\sum y \in s\ i\ \text{'space } M. \text{indicator } (s\ i\ -\ \{y\} \cap \text{space } M)\ x *_R y *_R cond\text{-exp } M\ F\ g\ x)$ **by** (intro Bochner-Integration.integral-sum[symmetric] integrable-mult-indicator[$OF - integrable\text{-}scaleR\text{-}right$]) auto
 also have $\dots = LINT\ x|M. s\ i\ x *_R cond\text{-exp } M\ F\ g\ x$ **using** $s\text{-scaleR-eq}$ **by** (intro Bochner-Integration.integral-cong) auto
 finally have $LINT\ x|M. s\ i\ x *_R g\ x = LINT\ x|?F. s\ i\ x *_R cond\text{-exp } M\ F\ g\ x$ **by** (simp add: integral-subalgebra2[$OF\ subalg$])
 }
 note $integral\text{-}s\text{-eq} = this$

show $integrable\ M\ (\lambda x. z\ x *_R cond\text{-exp } M\ F\ g\ x)$ **using** $s\text{-scaleR-cond-exp-g-meas}$
 asm(2) borel-measurable-cond-exp' **by** (intro integrable-from-subalg[$OF\ subalg$] integrable-cond-exp integrable-dominated-convergence[$OF - - - s\text{-scaleR-cond-exp-g-tendsto}$
 $s\text{-scaleR-cond-exp-g-AE-bdd}$] (auto intro: measurable-from-subalg[$OF\ subalg$] integrable-in-subalg measurable-in-subalg subalg))

have ($\lambda i. \text{LINT } x|M. s \ i \ x *_R g \ x$) \longrightarrow $\text{LINT } x|M. z \ x *_R g \ x$ **using**
s-scaleR-g-meas asm(1)[THEN integrable-norm] asm' borel-measurable-cond-exp'
by (*intro integral-dominated-convergence[OF - - - s-scaleR-g-tendsto s-scaleR-g-AE-bdd]*)
(auto intro: measurable-from-subalg[OF subalg])
moreover have ($\lambda i. \text{LINT } x|?F. s \ i \ x *_R \text{cond-exp } M \ F \ g \ x$) \longrightarrow
 $\text{LINT } x|?F. z \ x *_R \text{cond-exp } M \ F \ g \ x$ **using** *s-scaleR-cond-exp-g-meas asm(2)*
borel-measurable-cond-exp' **by** (*intro integral-dominated-convergence[OF - - - s-scaleR-cond-exp-g-tendsto*
s-scaleR-cond-exp-g-AE-bdd]) *(auto intro: measurable-from-subalg[OF subalg] inte-*
grable-in-subalg measurable-in-subalg subalg)
ultimately show $\text{LINT } x|M. z \ x *_R g \ x = \text{LINT } x|M. z \ x *_R \text{cond-exp}$
 $M \ F \ g \ x$ **using** *integral-s-eq using subalg* **by** (*simp add: LIMSEQ-unique inte-*
gral-subalgebra2)
qed
}
note $*$ = *this*

show *integrable* M ($\lambda x. f \ x *_R \text{cond-exp } M \ F \ g \ x$) **using** $*$ *assms measurable-*
in-subalg[OF subalg] **by** *blast*

{
fix A **assume** *asm: A ∈ F*
hence *integrable* M ($\lambda x. \text{indicat-real } A \ x *_R f \ x *_R g \ x$) **using** *subalg* **by**
(fastforce simp add: subalgebra-def intro!: integrable-mult-indicator assms(1))
hence *set-lebesgue-integral* $M \ A$ ($\lambda x. f \ x *_R g \ x$) = *set-lebesgue-integral* $M \ A$
 $(\lambda x. f \ x *_R \text{cond-exp } M \ F \ g \ x)$ **unfolding** *set-lebesgue-integral-def* **using** *asm* **by**
*(auto intro!: * measurable-in-subalg[OF subalg])*
}
thus *AE* x *in* $M. \text{cond-exp } M \ F$ ($\lambda x. f \ x *_R g \ x$) $x = f \ x *_R \text{cond-exp } M \ F \ g \ x$
using *borel-measurable-cond-exp* **by** (*intro cond-exp-charact, auto intro!: * assms*
measurable-in-subalg[OF subalg])
qed

lemma *cond-exp-sum [intro, simp]:*

fixes $f :: 't \Rightarrow 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes [*measurable*]: $\bigwedge i. \text{integrable } M \ (f \ i)$
shows *AE* x *in* $M. \text{cond-exp } M \ F$ ($\lambda x. \sum_{i \in I. f \ i \ x} x = (\sum_{i \in I. \text{cond-exp } M \ F}$
 $(f \ i) \ x)$
proof (*rule has-cond-exp-charact, intro has-cond-expI'*)
fix A **assume** [*measurable*]: $A \in \text{sets } F$
then have *A-meas* [*measurable*]: $A \in \text{sets } M$ **by** (*meson subsetD subalg subalge-*
bra-def)

have $(\int x \in A. (\sum_{i \in I. f \ i \ x}) \partial M) = (\int x. (\sum_{i \in I. \text{indicator } A \ x *_R f \ i \ x}) \partial M)$
unfolding *set-lebesgue-integral-def* **by** (*simp add: scaleR-sum-right*)
also have $\dots = (\sum_{i \in I. (\int x. \text{indicator } A \ x *_R f \ i \ x \ \partial M))$ **using** *assms* **by** (*auto*
intro!: Bochner-Integration.integral-sum integrable-mult-indicator)

also have $\dots = (\sum_{i \in I}. (\int x. \text{indicator } A \ x \ *_{\mathbb{R}} \text{cond-exp } M \ F \ (f \ i) \ x \ \partial M))$ **using** *cond-exp-set-integral[OF assms]* **by** (*simp add: set-lebesgue-integral-def*)
also have $\dots = (\int x. (\sum_{i \in I}. \text{indicator } A \ x \ *_{\mathbb{R}} \text{cond-exp } M \ F \ (f \ i) \ x) \partial M)$
using *assms* **by** (*auto intro!: Bochner-Integration.integral-sum[symmetric] integrable-mult-indicator*)
also have $\dots = (\int x \in A. (\sum_{i \in I}. \text{cond-exp } M \ F \ (f \ i) \ x) \partial M)$ **unfolding** *set-lebesgue-integral-def*
by (*simp add: scaleR-sum-right*)
finally show $(\int x \in A. (\sum_{i \in I}. f \ i \ x) \partial M) = (\int x \in A. (\sum_{i \in I}. \text{cond-exp } M \ F \ (f \ i) \ x) \partial M)$ **by** *auto*
qed (*auto simp add: assms integrable-cond-exp*)

8.1 Linearly Ordered Banach Spaces

lemma *cond-exp-gr-c*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes *integrable M f AE x in M. f x > c*

shows *AE x in M. cond-exp M F f x > c*

proof –

define X **where** $X = \{x \in \text{space } M. \text{cond-exp } M \ F \ f \ x \leq c\}$

have [*measurable*]: $X \in \text{sets } F$ **unfolding** X -*def* **by** *measurable (metis sets.top subalg subalgebra-def)*

hence X -*in-M*: $X \in \text{sets } M$ **using** *sets-restr-to-subalg subalg subalgebra-def* **by** *blast*

have *emeasure M X = 0*

proof (*rule ccontr*)

assume *emeasure M X ≠ 0*

have *emeasure (restr-to-subalg M F) X = emeasure M X* **by** (*simp add: emeasure-restr-to-subalg subalg*)

hence *emeasure (restr-to-subalg M F) X > 0* **using** $\neg(\text{emeasure } M \ X) = 0$ *gr-zeroI* **by** *auto*

then obtain A **where** $A: A \in \text{sets } (\text{restr-to-subalg } M \ F) \ A \subseteq X$ *emeasure (restr-to-subalg M F) A > 0* *emeasure (restr-to-subalg M F) A < ∞*

using *sigma-fin-subalg* **by** (*metis emeasure-notin-sets ennreal-0 infinity-ennreal-def le-less-linear neq-top-trans not-gr-zero order-refl sigma-finite-measure.approx-PInf-emeasure-with-finite*)

hence [*simp*]: $A \in \text{sets } F$ **using** *subalg sets-restr-to-subalg* **by** *blast*

hence A -*in-sets-M* [*simp*]: $A \in \text{sets } M$ **using** *sets-restr-to-subalg subalg subalgebra-def* **by** *blast*

have [*simp*]: *set-integrable M A* $(\lambda x. c)$ **using** A *subalg* **by** (*auto simp add: set-integrable-def emeasure-restr-to-subalg*)

have [*simp*]: *set-integrable M A f* **unfolding** *set-integrable-def* **by** (*rule integrable-mult-indicator, auto simp add: assms(1)*)

have *AE x in M. indicator A x *_ℝ c = indicator A x *_ℝ f x*

proof (*rule integral-eq-mono-AE-eq-AE*)

show *LINT x|M. indicator A x *_ℝ c = LINT x|M. indicator A x *_ℝ f x*

proof (*simp only: set-lebesgue-integral-def[symmetric], rule antisym*)

show $(\int x \in A. c \ \partial M) \leq (\int x \in A. f \ x \ \partial M)$ **using** *assms(2)* **by** (*intro set-integral-mono-AE-banach*) *auto*

have $(\int x \in A. f \ x \ \partial M) = (\int x \in A. \text{cond-exp } M \ F \ f \ x \ \partial M)$ **by** (*rule*

cond-exp-set-integral, auto simp add: assms)
also have $\dots \leq (\int x \in A. c \partial M)$ **using** A **by** (*auto intro!: set-integral-mono-banach*
simp add: X-def)
finally show $(\int x \in A. f x \partial M) \leq (\int x \in A. c \partial M)$ **by** *simp*
qed
show $AE x \text{ in } M. \text{indicator } A x *_R c \leq \text{indicator } A x *_R f x$ **using** *assms* **by**
(*auto simp add: X-def indicator-def*)
qed (*auto simp add: set-integrable-def[symmetric]*)
hence $AE x \in A \text{ in } M. c = f x$ **by** *auto*
hence $AE x \in A \text{ in } M. \text{False}$ **using** *assms(2)* **by** *auto*
hence $A \in \text{null-sets } M$ **using** *AE-iff-null-sets A-in-sets-M* **by** *metis*
thus False **using** $A(3)$ **by** (*simp add: emeasure-restr-to-subalg null-setsD1*
subalg)
qed
thus *?thesis* **using** *AE-iff-null-sets[OF X-in-M]* **unfolding** *X-def* **by** *auto*
qed

corollary *cond-exp-less-c:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$
dered-real-vector}
assumes *integrable M f AE x in M. f x < c*
shows $AE x \text{ in } M. \text{cond-exp } M F f x < c$
proof –
have $AE x \text{ in } M. \text{cond-exp } M F f x = - \text{cond-exp } M F (\lambda x. - f x) x$ **using**
cond-exp-uminus[OF assms(1)] **by** *auto*
moreover have $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. - f x) x > - c$ **using** *assms*
by (*intro cond-exp-gr-c*) *auto*
ultimately show *?thesis* **by** (*force simp add: minus-less-iff*)
qed

lemma *cond-exp-mono-strict:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$
dered-real-vector}
assumes *integrable M f integrable M g AE x in M. f x < g x*
shows $AE x \text{ in } M. \text{cond-exp } M F f x < \text{cond-exp } M F g x$
using *cond-exp-less-c[OF Bochner-Integration.integrable-diff, OF assms(1,2), of*
0]
cond-exp-diff[OF assms(1,2)] assms(3) **by** *auto*

lemma *cond-exp-ge-c:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$
dered-real-vector}
assumes [*measurable*]: *integrable M f*
and $AE x \text{ in } M. f x \geq c$
shows $AE x \text{ in } M. \text{cond-exp } M F f x \geq c$
proof –
let $?F = \text{restr-to-subalg } M F$
interpret *sigma-finite-measure restr-to-subalg M F* **using** *sigma-fin-subalg* **by**
auto

```

{
  fix A assume asm: A ∈ sets ?F 0 < measure ?F A
  have [simp]: sets ?F = sets F measure ?F A = measure M A using asm by (auto
simp add: measure-def sets-restr-to-subalg[OF subalg] emeasure-restr-to-subalg[OF
subalg])
  have M-A: emeasure M A < ∞ using measure-zero-top asm by (force simp
add: top.not-eq-extremum)
  hence F-A: emeasure ?F A < ∞ using asm(1) emeasure-restr-to-subalg subalg
by fastforce
  have set-lebesgue-integral M A (λ-. c) ≤ set-lebesgue-integral M A f using
assms asm M-A subalg by (intro set-integral-mono-AE-banach, auto simp add:
set-integrable-def integrable-mult-indicator subalgebra-def sets-restr-to-subalg)
  also have ... = set-lebesgue-integral M A (cond-exp M F f) using cond-exp-set-integral[OF
assms(1)] asm by auto
  also have ... = set-lebesgue-integral ?F A (cond-exp M F f) unfolding set-lebesgue-integral-def
using asm borel-measurable-cond-exp by (intro integral-subalgebra2[OF subalg, sym-
metric], simp)
  finally have (1 / measure ?F A) *R set-lebesgue-integral ?F A (cond-exp M F f)
∈ {c..} using asm subalg M-A by (auto simp add: set-integral-const subalgebra-def
intro!: pos-divideR-le-eq[THEN iffD1])
}
thus ?thesis using AE-restr-to-subalg[OF subalg] averaging-theorem[OF inte-
grable-in-subalg closed-atLeast, OF subalg borel-measurable-cond-exp integrable-cond-exp]
by auto
qed

```

corollary *cond-exp-le-c*:

```

fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
dered-real-vector}
assumes integrable M f
and AE x in M. f x ≤ c
shows AE x in M. cond-exp M F f x ≤ c
proof -
  have AE x in M. cond-exp M F f x = - cond-exp M F (λx. - f x) x using
cond-exp-uminus[OF assms(1)] by force
  moreover have AE x in M. cond-exp M F (λx. - f x) x ≥ - c using assms
by (intro cond-exp-ge-c) auto
  ultimately show ?thesis by (force simp add: minus-le-iff)
qed

```

corollary *cond-exp-mono*:

```

fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
dered-real-vector}
assumes integrable M f integrable M g AE x in M. f x ≤ g x
shows AE x in M. cond-exp M F f x ≤ cond-exp M F g x
using cond-exp-le-c[OF Bochner-Integration.integrable-diff, OF assms(1,2), of
0]
cond-exp-diff[OF assms(1,2)] assms(3) by auto

```

corollary *cond-exp-min*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes *integrable* $M f$ *integrable* $M g$

shows $AE \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \min (f x) (g x)) \xi \leq \min (\text{cond-exp } M F f \xi) (\text{cond-exp } M F g \xi)$

proof –

have $AE \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \min (f x) (g x)) \xi \leq \text{cond-exp } M F f \xi$ **by** (*intro cond-exp-mono integrable-min assms, simp*)

moreover have $AE \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \min (f x) (g x)) \xi \leq \text{cond-exp } M F g \xi$ **by** (*intro cond-exp-mono integrable-min assms, simp*)

ultimately show $AE \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \min (f x) (g x)) \xi \leq \min (\text{cond-exp } M F f \xi) (\text{cond-exp } M F g \xi)$ **by** *fastforce*

qed

corollary *cond-exp-max*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes *integrable* $M f$ *integrable* $M g$

shows $AE \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \max (f x) (g x)) \xi \geq \max (\text{cond-exp } M F f \xi) (\text{cond-exp } M F g \xi)$

proof –

have $AE \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \max (f x) (g x)) \xi \geq \text{cond-exp } M F f \xi$ **by** (*intro cond-exp-mono integrable-max assms, simp*)

moreover have $AE \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \max (f x) (g x)) \xi \geq \text{cond-exp } M F g \xi$ **by** (*intro cond-exp-mono integrable-max assms, simp*)

ultimately show $AE \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \max (f x) (g x)) \xi \geq \max (\text{cond-exp } M F f \xi) (\text{cond-exp } M F g \xi)$ **by** *fastforce*

qed

corollary *cond-exp-inf*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector, lattice}\}$

assumes *integrable* $M f$ *integrable* $M g$

shows $AE \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \inf (f x) (g x)) \xi \leq \inf (\text{cond-exp } M F f \xi) (\text{cond-exp } M F g \xi)$

unfolding *inf-min using assms by (rule cond-exp-min)*

corollary *cond-exp-sup*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector, lattice}\}$

assumes *integrable* $M f$ *integrable* $M g$

shows $AE \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \sup (f x) (g x)) \xi \geq \sup (\text{cond-exp } M F f \xi) (\text{cond-exp } M F g \xi)$

unfolding *sup-max using assms by (rule cond-exp-max)*

end

8.2 Probability Spaces

lemma (in *prob-space*) *prob-space-restr-to-subalg*:

assumes *subalgebra* $M\ F$

shows *prob-space* (*restr-to-subalg* $M\ F$)

proof –

have *countable* {*space* M } **by** *simp*

moreover have {*space* M } \subseteq *sets* (*restr-to-subalg* $M\ F$) **unfolding** *restr-to-subalg-def*

by *simp*

moreover have \bigcup {*space* M } = *space* (*restr-to-subalg* $M\ F$) **unfolding** *space-restr-to-subalg*

by *simp*

moreover have $\forall a \in \{\text{space } M\}. \text{emeasure } (\text{restr-to-subalg } M\ F) \ a \neq \infty$ **unfolding** *restr-to-subalg-def* *emeasure-measure-of-conv* **by** *simp*

ultimately show *prob-space* (*restr-to-subalg* $M\ F$) **using** *emeasure-space-1* *emeasure-restr-to-subalg* [*OF* *assms* *sets.top*] *assms*

by *unfold-locale* (*blast*, *auto* *simp* *add*: *space-restr-to-subalg* *subalgebra-def*)

qed

lemma (in *prob-space*) *sigma-finite-subalgebra-restr-to-subalg*:

assumes *subalgebra* $M\ F$

shows *sigma-finite-subalgebra* $M\ F$

proof (*intro* *sigma-finite-subalgebra.intro*)

interpret F : *prob-space* *restr-to-subalg* $M\ F$ **using** *assms* *prob-space-restr-to-subalg*

by *blast*

show *sigma-finite-measure* (*restr-to-subalg* $M\ F$) **by** (*rule* F .*sigma-finite-measure-axioms*)

qed (*rule* *assms*)

lemma (in *prob-space*) *cond-exp-trivial*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$

assumes *integrable* $M\ f$

shows $\text{AE } x \text{ in } M. \text{cond-exp } M \ (\text{sigma } (\text{space } M) \ \{\}) \ f \ x = \text{expectation } f$

proof –

interpret *sigma-finite-subalgebra* M *sigma* (*space* M) { $\}$ **by** (*auto* *intro*: *sigma-finite-subalgebra-restr-to-subalg* *simp* *add*: *subalgebra-def* *sigma-sets-empty-eq*)

show *?thesis* **using** *assms* **by** (*intro* *cond-exp-charact*) (*auto* *simp* *add*: *sigma-sets-empty-eq* *set-lebesgue-integral-def* *prob-space* *cong*: *Bochner-Integration.integral-cong*)

qed

lemma (in *prob-space*) *cond-exp-indep-subalgebra*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}, \text{real-normed-field}\}$

assumes *subalgebra*: *subalgebra* $M\ F$ *subalgebra* $M\ G$

and *independent*: *indep-set* G (*sigma* (*space* M) ($F \cup \text{vimage-algebra } (\text{space } M) \ f \ \text{borel}$))

assumes [*measurable*]: *integrable* $M\ f$

shows $\text{AE } x \text{ in } M. \text{cond-exp } M \ (\text{sigma } (\text{space } M) \ (F \cup G)) \ f \ x = \text{cond-exp } M\ F \ f \ x$

proof –

interpret *Un-sigma*: *sigma-finite-subalgebra* M *sigma* (*space* M) ($F \cup G$) **using** *assms*(1,2) **by** (*auto* *intro*!: *sigma-finite-subalgebra-restr-to-subalg* *sets.sigma-sets-subset* *simp* *add*: *subalgebra-def* *space-measure-of-conv* *sets-measure-of-conv*)


```

interpret sigma-finite-subalgebra  $M$   $F$  using assms by (auto intro: sigma-finite-subalgebra-restr-to-subalg)
{
  fix  $A$ 
  assume asm:  $A \in \text{sigma } (\text{space } M) \{a \cap b \mid a \text{ b. } a \in F \wedge b \in G\}$ 
  have in-events: sigma-sets (space  $M$ )  $\{a \cap b \mid a \text{ b. } a \in \text{sets } F \wedge b \in \text{sets } G\} \subseteq \text{events}$  using subalgebra by (intro sets.sigma-sets-subset, auto simp add: subalgebra-def)
  have Int-stable  $\{a \cap b \mid a \text{ b. } a \in F \wedge b \in G\}$ 
  proof (intro Int-stableI, clarsimp)
    fix  $af \text{ } bf \text{ } ag \text{ } bg$ 
    assume  $F$ :  $af \in F \text{ } bf \in F$  and  $G$ :  $ag \in G \text{ } bg \in G$ 
    have  $af \cap bf \in F$  by (intro sets.Int F)
    moreover have  $ag \cap bg \in G$  by (intro sets.Int G)
    ultimately show  $\exists a \text{ b. } af \cap ag \cap (bf \cap bg) = a \cap b \wedge a \in \text{sets } F \wedge b \in \text{sets } G$  by (metis inf-assoc inf-left-commute)
  qed
  moreover have  $\{a \cap b \mid a \text{ b. } a \in F \wedge b \in G\} \subseteq \text{Pow } (\text{space } M)$  using subalgebra by (force simp add: subalgebra-def dest: sets.sets-into-space)
  moreover have  $A \in \text{sigma-sets } (\text{space } M) \{a \cap b \mid a \text{ b. } a \in F \wedge b \in G\}$  using calculation asm by force
  ultimately have set-lebesgue-integral  $M$   $A$   $f = \text{set-lebesgue-integral } M$   $A$  (cond-exp  $M$   $F$   $f$ )
  proof (induction rule: sigma-sets-induct-disjoint)
    case (basic  $A$ )
    then obtain  $a \text{ } b$  where  $A$ :  $A = a \cap b \text{ } a \in F \text{ } b \in G$  by blast

    hence events[measurable]:  $a \in \text{events } b \in \text{events}$  using subalgebra by (auto simp add: subalgebra-def)

    have [simp]: sigma-sets (space  $M$ )  $\{\text{indicator } b - 'A \cap \text{space } M \mid A. A \in \text{borel}\} \subseteq G$ 
    using borel-measurable-indicator[OF  $A(3)$ , THEN measurable-sets] sets.top subalgebra
    by (intro sets.sigma-sets-subset') (fastforce simp add: subalgebra-def)+

    have Un-in-sigma:  $F \cup \text{vimage-algebra } (\text{space } M) \text{ } f \text{ borel} \subseteq \text{sigma } (\text{space } M) (F \cup \text{vimage-algebra } (\text{space } M) \text{ } f \text{ borel})$  by (metis equalityE le-supI sets.space-closed sigma-le-sets space-vimage-algebra subalg subalgebra-def)

    have [intro]: indep-var borel (indicator  $b$ ) borel  $(\lambda \omega. \text{indicator } a \text{ } \omega *_{\mathbb{R}} f \text{ } \omega)$ 
    proof –
      have [simp]: sigma-sets (space  $M$ )  $\{(\lambda \omega. \text{indicator } a \text{ } \omega *_{\mathbb{R}} f \text{ } \omega) - 'A \cap \text{space } M \mid A. A \in \text{borel}\} \subseteq \text{sigma } (\text{space } M) (F \cup \text{vimage-algebra } (\text{space } M) \text{ } f \text{ borel})$ 
      proof –
        have  $*$ :  $(\lambda \omega. \text{indicator } a \text{ } \omega *_{\mathbb{R}} f \text{ } \omega) \in \text{borel-measurable } (\text{sigma } (\text{space } M) (F \cup \text{vimage-algebra } (\text{space } M) \text{ } f \text{ borel}))$ 
        using borel-measurable-indicator[OF  $A(2)$ , THEN measurable-sets, OF borel-open] subalgebra
        by (intro borel-measurable-scaleR borel-measurableI Un-in-sigma [THEN

```

```

subsetD])
  (auto simp add: space-measure-of-conv subalgebra-def sets-vimage-algebra2)
  thus ?thesis using measurable-sets[OF *] by (intro sets.sigma-sets-subset',
auto simp add: space-measure-of-conv)
qed
  have indep-set (sigma-sets (space M) {indicator b - ' A ∩ space M | A. A ∈
borel}) (sigma-sets (space M) {(λω. indicator a ω *R f ω) - ' A ∩ space M | A. A
∈ borel})
  using independent unfolding indep-set-def by (rule indep-sets-mono-sets,
auto split: bool.split)
  thus ?thesis by (subst indep-var-eq, auto intro!: borel-measurable-scaleR)
qed

  have [intro]: indep-var borel (indicator b) borel (λω. indicat-real a ω *R
cond-exp M F f ω)
  proof -
    have [simp]: sigma-sets (space M) {(λω. indicator a ω *R cond-exp M F f
ω) - ' A ∩ space M | A. A ∈ borel} ⊆ sigma (space M) (F ∪ vimage-algebra (space
M) f borel)
    proof -
      have *: (λω. indicator a ω *R cond-exp M F f ω) ∈ borel-measurable (sigma
(space M) (F ∪ vimage-algebra (space M) f borel))
      using borel-measurable-indicator[OF A(2), THEN measurable-sets, OF
borel-open] subalgebra
        borel-measurable-cond-exp[THEN measurable-sets, OF borel-open, of
- M F f]
      by (intro borel-measurable-scaleR borel-measurableI Un-in-sigma[THEN
subsetD])
    (auto simp add: space-measure-of-conv subalgebra-def)
  thus ?thesis using measurable-sets[OF *] by (intro sets.sigma-sets-subset',
auto simp add: space-measure-of-conv)
qed
  have indep-set (sigma-sets (space M) {indicator b - ' A ∩ space M | A. A ∈
borel}) (sigma-sets (space M) {(λω. indicator a ω *R cond-exp M F f ω) - ' A ∩
space M | A. A ∈ borel})
  using independent unfolding indep-set-def by (rule indep-sets-mono-sets,
auto split: bool.split)
  thus ?thesis by (subst indep-var-eq, auto intro!: borel-measurable-scaleR)
qed

  have set-lebesgue-integral M A f = (LINT x|M. indicator b x * (indicator a
x *R f x))
  unfolding set-lebesgue-integral-def A indicator-inter-arith
  by (intro Bochner-Integration.integral-cong, auto simp add: scaleR-scaleR[symmetric]
indicator-times-eq-if(1))
  also have ... = (LINT x|M. indicator b x) * (LINT x|M. indicator a x *R f
x)
  by (intro indep-var-lebesgue-integral

```

```

      Bochner-Integration.integrable-bound[OF integrable-const[of 1 :: 'b]
borel-measurable-indicator]
      integrable-mult-indicator[OF - assms(4)], blast) (auto simp add:
indicator-def)
    also have ... = (LINT x|M. indicator b x) * (LINT x|M. indicator a x *R
cond-exp M F f x)
    using cond-exp-set-integral[OF assms(4) A(2)] unfolding set-lebesgue-integral-def
by argo
    also have ... = (LINT x|M. indicator b x * (indicator a x *R cond-exp M
F f x))
    by (intro indep-var-lebesgue-integral[symmetric]
      Bochner-Integration.integrable-bound[OF integrable-const[of 1 :: 'b]
borel-measurable-indicator]
      integrable-mult-indicator[OF - integrable-cond-exp], blast) (auto simp
add: indicator-def)
    also have ... = set-lebesgue-integral M A (cond-exp M F f)
    unfolding set-lebesgue-integral-def A indicator-inter-arith
    by (intro Bochner-Integration.integral-cong, auto simp add: scaleR-scaleR[symmetric]
indicator-times-eq-if(1))
    finally show ?case .
  next
    case empty
    then show ?case unfolding set-lebesgue-integral-def by simp
  next
    case (compl A)
    have A-in-space: A ⊆ space M using compl using in-events sets.sets-into-space
by blast
    have set-lebesgue-integral M (space M - A) f = set-lebesgue-integral M (space
M - A ∪ A) f - set-lebesgue-integral M A f
    using compl(1) in-events
    by (subst set-integral-Un[of space M - A A], blast)
      (simp | intro integrable-mult-indicator[folded set-integrable-def, OF -
assms(4)], fast)+
    also have ... = set-lebesgue-integral M (space M - A ∪ A) (cond-exp M F f)
- set-lebesgue-integral M A (cond-exp M F f)
    using cond-exp-set-integral[OF assms(4) sets.top] compl subalgebra by (simp
add: subalgebra-def Un-absorb2[OF A-in-space])
    also have ... = set-lebesgue-integral M (space M - A) (cond-exp M F f)
    using compl(1) in-events
    by (subst set-integral-Un[of space M - A A], blast)
      (simp | intro integrable-mult-indicator[folded set-integrable-def, OF -
integrable-cond-exp], fast)+
    finally show ?case .
  next
    case (union A)
    have set-lebesgue-integral M (⋃ (range A)) f = (∑ i. set-lebesgue-integral M
(A i) f)
    using union in-events
    by (intro lebesgue-integral-countable-add) (auto simp add: disjoint-family-onD

```

```

intro!: integrable-mult-indicator[folded set-integrable-def, OF - assms(4)]
  also have ... = ( $\sum i$ . set-lebesgue-integral M (A i) (cond-exp M F f)) using
union by presburger
  also have ... = set-lebesgue-integral M ( $\bigcup$  (range A)) (cond-exp M F f)
  using union in-events
  by (intro lebesgue-integral-countable-add[symmetric]) (auto simp add: dis-
joint-family-onD intro!: integrable-mult-indicator[folded set-integrable-def, OF - in-
tegrable-cond-exp])
  finally show ?case .
qed
}
moreover have sigma (space M) {a  $\cap$  b | a b. a  $\in$  F  $\wedge$  b  $\in$  G} = sigma (space
M) (F  $\cup$  G)
proof -
  have sigma-sets (space M) {a  $\cap$  b | a b. a  $\in$  sets F  $\wedge$  b  $\in$  sets G} = sigma-sets
(space M) (sets F  $\cup$  sets G)
  proof (intro sigma-sets-eqI ; clarsimp ; cases)
    fix a b assume asm: a  $\in$  F b  $\in$  G
    thus a  $\cap$  b  $\in$  sigma-sets (space M) (F  $\cup$  G) using subalgebra unfolding
Int-range-binary by (intro sigma-sets-Inter[OF - binary-in-sigma-sets]) (force simp
add: subalgebra-def dest: sets.sets-into-space)+
  next
    fix a
    assume a  $\in$  sets F
    thus a  $\in$  sigma-sets (space M) {a  $\cap$  b | a b. a  $\in$  sets F  $\wedge$  b  $\in$  sets G}
      using subalgebra sets.top[of G] sets.sets-into-space[of - F]
      by (intro sigma-sets.Basic, auto simp add: subalgebra-def)
  next
    fix a assume a  $\in$  sets F  $\vee$  a  $\in$  sets G a  $\notin$  sets F
    hence a  $\in$  sets G by blast
    thus a  $\in$  sigma-sets (space M) {a  $\cap$  b | a b. a  $\in$  sets F  $\wedge$  b  $\in$  sets G}
      using subalgebra sets.top[of F] sets.sets-into-space[of - G]
      by (intro sigma-sets.Basic, auto simp add: subalgebra-def)
    qed (blast)
    thus ?thesis using subalgebra by (intro sigma-eqI) (force simp add: subalge-
bra-def dest: sets.sets-into-space)+
  qed
moreover have (cond-exp M F f)  $\in$  borel-measurable (sigma (space M) (sets F
 $\cup$  sets G))
proof -
  have F  $\subseteq$  sigma (space M) (F  $\cup$  G) by (metis Un-least Un-upper1 mea-
sure-of-of-measure sets.space-closed sets-measure-of sigma-sets-subseteq subalg sub-
algebra(2) subalgebra-def)
  thus ?thesis using borel-measurable-cond-exp[THEN measurable-sets, OF borel-open,
of - M F f] subalgebra by (intro borel-measurableI, force simp only: space-measure-of-conv
subalgebra-def)
  qed
ultimately show ?thesis using assms(4) integrable-cond-exp by (intro Un-sigma.cond-exp-charact)
presburger+

```

qed

lemma (in prob-space) cond-exp-indep:
 fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, real-normed-field}\}$
 assumes subalgebra: subalgebra M F
 and independent: indep-set F (vimage-algebra (space M) f borel)
 and integrable: integrable M f
 shows AE x in M . cond-exp M F f $x = \text{expectation } f$
proof –
 have indep-set F (sigma (space M) (sigma (space M) $\{\}$ \cup (vimage-algebra (space M) f borel)))
 using independent unfolding indep-set-def
 by (rule indep-sets-mono-sets, simp add: bool.split)
 (metis bot.extremum dual-order.refl sets.sets-measure-of-eq sets.sigma-sets-subset'
 sets-vimage-algebra-space space-vimage-algebra sup.absorb-iff2)
 hence cond-exp-indep: AE x in M . cond-exp M (sigma (space M) (sigma (space M) $\{\}$ \cup F)) f $x = \text{expectation } f$
 using cond-exp-indep-subalgebra[OF - subalgebra - integrable, of sigma (space M) $\{\}$] cond-exp-trivial[OF integrable]
 by (auto simp add: subalgebra-def sigma-sets-empty-eq)
 have sets (sigma (space M) (sigma (space M) $\{\}$ \cup F)) = F
 using subalgebra sets.top[of F] unfolding subalgebra-def
 by (simp add: sigma-sets-empty-eq, subst insert-absorb[of space M F], blast)
 (metis insert-absorb[OF sets.empty-sets] sets.sets-measure-of-eq)
 hence AE x in M . cond-exp M (sigma (space M) (sigma (space M) $\{\}$ \cup F)) f
 $x = \text{cond-exp } M$ F f x by (rule cond-exp-sets-cong)
 thus ?thesis using cond-exp-indep by force
 qed
 end

theory Filtered-Measure
 imports HOL-Probability.Conditional-Expectation
 begin

9 Filtered Measure Spaces

9.1 Filtered Measure

locale filtered-measure =
 fixes M F and $t_0 :: 'b :: \{\text{second-countable-topology, order-topology, t2-space}\}$
 assumes subalgebra: $\bigwedge i. t_0 \leq i \implies \text{subalgebra } M$ (F i)
 and sets- F -mono: $\bigwedge i$ $j. t_0 \leq i \implies i \leq j \implies \text{sets } (F$ $i) \leq \text{sets } (F$ $j)$
 begin

lemma space- F [simp]:
 assumes $t_0 \leq i$
 shows space (F i) = space M

```

using subalgebra assms by (simp add: subalgebra-def)

lemma subalgebra-F[intro]:
  assumes  $t_0 \leq i \leq j$ 
  shows subalgebra (F j) (F i)
  unfolding subalgebra-def using assms by (simp add: sets-F-mono)

lemma borel-measurable-mono:
  assumes  $t_0 \leq i \leq j$ 
  shows borel-measurable (F i)  $\subseteq$  borel-measurable (F j)
  unfolding subset-iff by (metis assms subalgebra-F measurable-from-subalg)

end

locale linearly-filtered-measure = filtered-measure M F t_0 for M and F :: - ::
{linorder-topology}  $\Rightarrow$  - and t_0

locale nat-filtered-measure = linearly-filtered-measure M F 0 for M and F :: nat
 $\Rightarrow$  -
locale real-filtered-measure = linearly-filtered-measure M F 0 for M and F :: real
 $\Rightarrow$  -

```

9.2 Sigma Finite Filtered Measure

The locale presented here is a generalization of the *sigma-finite-subalgebra* for a particular filtration.

```

locale sigma-finite-filtered-measure = filtered-measure +
  assumes sigma-finite-initial: sigma-finite-subalgebra M (F t_0)

lemma (in sigma-finite-filtered-measure) sigma-finite-subalgebra-F[intro]:
  assumes  $t_0 \leq i$ 
  shows sigma-finite-subalgebra M (F i)
  using assms by (metis dual-order.refl sets-F-mono sigma-finite-initial sigma-finite-subalgebra.nested-subalg-is
subalgebra subalgebra-def)

locale nat-sigma-finite-filtered-measure = sigma-finite-filtered-measure M F 0 ::
nat for M F
locale real-sigma-finite-filtered-measure = sigma-finite-filtered-measure M F 0 ::
real for M F

sublocale nat-sigma-finite-filtered-measure  $\subseteq$  sigma-finite-subalgebra M F i by
blast
sublocale real-sigma-finite-filtered-measure  $\subseteq$  sigma-finite-subalgebra M F |i| by
fastforce

```

9.3 Finite Filtered Measure

```

locale finite-filtered-measure = filtered-measure + finite-measure

```

```

sublocale finite-filtered-measure  $\subseteq$  sigma-finite-filtered-measure
  using subalgebra by (unfold-locales, blast, meson dual-order.refl finite-measure-axioms
finite-measure-def finite-measure-restr-to-subalg sigma-finite-measure.sigma-finite-countable
subalgebra)

```

```

locale nat-finite-filtered-measure = finite-filtered-measure M F 0 :: nat for M F
locale real-finite-filtered-measure = finite-filtered-measure M F 0 :: real for M F

```

```

sublocale nat-finite-filtered-measure  $\subseteq$  nat-sigma-finite-filtered-measure ..
sublocale real-finite-filtered-measure  $\subseteq$  real-sigma-finite-filtered-measure ..

```

9.4 Constant Filtration

```

lemma filtered-measure-constant-filtration:
  assumes subalgebra M F
  shows filtered-measure M ( $\lambda$ -. F) t0
  using assms by (unfold-locales) blast

```

```

sublocale sigma-finite-subalgebra  $\subseteq$  constant-filtration: sigma-finite-filtered-measure
M  $\lambda$ -. :: 't :: {second-countable-topology, linorder-topology}. F t0
  using subalg by (unfold-locales) blast

```

```

lemma (in finite-measure) filtered-measure-constant-filtration:
  assumes subalgebra M F
  shows finite-filtered-measure M ( $\lambda$ -. F) t0
  using assms by (unfold-locales) blast

```

end

```

theory Stochastic-Process
imports Filtered-Measure Measure-Space-Supplement
begin

```

10 Stochastic Processes

10.1 Stochastic Process

A stochastic process is a collection of random variables, indexed by a type '*b*.

```

locale stochastic-process =
  fixes M t0 and X :: 'b :: {second-countable-topology, order-topology, t2-space}  $\Rightarrow$ 
'a  $\Rightarrow$  'c :: {second-countable-topology, banach}
  assumes random-variable[measurable]:  $\bigwedge i. t_0 \leq i \Rightarrow X\ i \in \text{borel-measurable } M$ 
begin

```

```

definition left-continuous where left-continuous = (AE  $\xi$  in M.  $\forall t. \text{continuous}$ 
(at-left t) ( $\lambda i. X\ i\ \xi$ ))

```

definition *right-continuous* **where** *right-continuous* = ($\text{AE } \xi \text{ in } M. \forall t. \text{continuous} \text{ (at-right } t) (\lambda i. X \ i \ \xi)$)

end

locale *nat-stochastic-process* = *stochastic-process* $M \ 0 :: \text{nat } X$ **for** $M \ X$
locale *real-stochastic-process* = *stochastic-process* $M \ 0 :: \text{real } X$ **for** $M \ X$

lemma *stochastic-process-const-fun*:
assumes $f \in \text{borel-measurable } M$
shows *stochastic-process* $M \ t_0 (\lambda -. f)$ **using** *assms* **by** (*unfold-locale*)

lemma *stochastic-process-const*:
shows *stochastic-process* $M \ t_0 (\lambda i -. c \ i)$ **by** (*unfold-locale*) *simp*

context *stochastic-process*
begin

lemma *compose-stochastic*:
assumes $\bigwedge i. t_0 \leq i \implies f \ i \in \text{borel-measurable borel}$
shows *stochastic-process* $M \ t_0 (\lambda i \xi. (f \ i) (X \ i \ \xi))$
by (*unfold-locale*) (*intro measurable-compose[OF random-variable assms]*)

lemma *norm-stochastic*: *stochastic-process* $M \ t_0 (\lambda i \xi. \text{norm } (X \ i \ \xi))$ **by** (*fastforce intro: compose-stochastic*)

lemma *scaleR-right-stochastic*:
assumes *stochastic-process* $M \ t_0 \ Y$
shows *stochastic-process* $M \ t_0 (\lambda i \xi. (Y \ i \ \xi) *_R (X \ i \ \xi))$
using *stochastic-process.random-variable[OF assms]* *random-variable* **by** (*unfold-locale*) *simp*

lemma *scaleR-right-const-fun-stochastic*:
assumes $f \in \text{borel-measurable } M$
shows *stochastic-process* $M \ t_0 (\lambda i \xi. f \ \xi *_R (X \ i \ \xi))$
by (*unfold-locale*) (*intro borel-measurable-scaleR assms random-variable*)

lemma *scaleR-right-const-stochastic*: *stochastic-process* $M \ t_0 (\lambda i \xi. c \ i *_R (X \ i \ \xi))$
by (*unfold-locale*) *simp*

lemma *add-stochastic*:
assumes *stochastic-process* $M \ t_0 \ Y$
shows *stochastic-process* $M \ t_0 (\lambda i \xi. X \ i \ \xi + Y \ i \ \xi)$
using *stochastic-process.random-variable[OF assms]* *random-variable* **by** (*unfold-locale*) *simp*

lemma *diff-stochastic*:
assumes *stochastic-process* $M \ t_0 \ Y$
shows *stochastic-process* $M \ t_0 (\lambda i \xi. X \ i \ \xi - Y \ i \ \xi)$

using *stochastic-process.random-variable*[*OF assms*] *random-variable* **by** (*unfold-locales*)
simp

lemma *uminus-stochastic: stochastic-process* $M\ t_0\ (-X)$ **using** *scaleR-right-const-stochastic*[*of*
 $\lambda\cdot, -1]$ **by** (*simp add: fun-Compl-def*)

lemma *partial-sum-stochastic: stochastic-process* $M\ t_0\ (\lambda n\ \xi. \sum_{i \in \{t_0..n\}} X\ i\ \xi)$
by (*unfold-locales*) *simp*

lemma *partial-sum'-stochastic: stochastic-process* $M\ t_0\ (\lambda n\ \xi. \sum_{i \in \{t_0..<n\}} X\ i\ \xi)$
by (*unfold-locales*) *simp*

end

lemma *stochastic-process-sum:*

assumes $\bigwedge i. i \in I \implies \text{stochastic-process } M\ t_0\ (X\ i)$
shows *stochastic-process* $M\ t_0\ (\lambda k\ \xi. \sum_{i \in I} X\ i\ k\ \xi)$ **using** *assms*[*THEN*
stochastic-process.random-variable] **by** (*unfold-locales, auto*)

10.1.1 Natural Filtration

The natural filtration induced by a stochastic process X is the filtration generated by all events involving the process up to the time index t , i.e. $\Sigma\ t = \sigma\ \{X\ s \mid s. s \leq t\}$.

definition *natural-filtration* :: '*a measure* \Rightarrow '*b* \Rightarrow ('*b* \Rightarrow '*a* \Rightarrow '*c* :: *topological-space*) \Rightarrow '*b* :: {*second-countable-topology, order-topology*} \Rightarrow '*a measure* **where**
natural-filtration $M\ t_0\ Y = (\lambda t. \text{family-vimage-algebra } (\text{space } M) \{Y\ i \mid i. i \in \{t_0..t\}\})\ \text{borel}$

abbreviation *nat-natural-filtration* $\equiv \lambda M. \text{natural-filtration } M\ (0 :: \text{nat})$

abbreviation *real-natural-filtration* $\equiv \lambda M. \text{natural-filtration } M\ (0 :: \text{real})$

lemma *space-natural-filtration*[*simp*]: *space* (*natural-filtration* $M\ t_0\ X\ t$) = *space*
 M **unfolding** *natural-filtration-def* *space-family-vimage-algebra* ..

lemma *sets-natural-filtration: sets* (*natural-filtration* $M\ t_0\ X\ t$) = *sigma-sets* (*space*
 M) $(\bigcup_{i \in \{t_0..t\}} \{X\ i - 'A \cap \text{space } M \mid A. A \in \text{borel}\})$

unfolding *natural-filtration-def* *sets-family-vimage-algebra* **by** (*intro sigma-sets-eqI*)
blast+

lemma *sets-natural-filtration'*:

assumes *borel* = *sigma UNIV S*

shows *sets* (*natural-filtration* $M\ t_0\ X\ t$) = *sigma-sets* (*space* M) $(\bigcup_{i \in \{t_0..t\}} \{X\ i - 'A \cap \text{space } M \mid A. A \in S\})$

proof (*subst sets-natural-filtration, intro sigma-sets-eqI, clarify*)

fix i **and** $A :: 'a \text{ set}$ **assume** *asm*: $i \in \{t_0..t\}\ A \in \text{sets borel}$

hence $A \in \text{sigma-sets UNIV } S$ **unfolding** *assms* **by** *simp*

thus $X\ i - 'A \cap \text{space } M \in \text{sigma-sets } (\text{space } M) (\bigcup_{i \in \{t_0..t\}} \{X\ i - 'A \cap \text{space } M \mid A. A \in S\})$

```

proof (induction)
  case (Compl a)
    have  $X\ i - ' (UNIV - a) \cap \text{space } M = \text{space } M - (X\ i - ' a \cap \text{space } M)$  by
    blast
    then show ?case using Compl(2)[THEN sigma-sets.Compl] by presburger
  next
    case (Union a)
    have  $X\ i - ' \bigcup (\text{range } a) \cap \text{space } M = \bigcup (\text{range } (\lambda j. X\ i - ' a\ j \cap \text{space } M))$ 
by blast
    then show ?case using Union(2)[THEN sigma-sets.Union] by presburger
  qed (auto intro: asm sigma-sets.Empty)
qed (intro sigma-sets.Basic, force simp add: assms)

```

lemma sets-natural-filtration-open:

```

  sets (natural-filtration  $M\ t_0\ X\ t$ ) = sigma-sets (space  $M$ ) ( $\bigcup_{i \in \{t_0..t\}} \{X\ i - ' A \cap \text{space } M \mid A. \text{open } A\}$ )
using sets-natural-filtration' by (force simp only: borel-def mem-Collect-eq)

```

lemma sets-natural-filtration-oi:

```

  sets (natural-filtration  $M\ t_0\ X\ t$ ) = sigma-sets (space  $M$ ) ( $\bigcup_{i \in \{t_0..t\}} \{X\ i - ' A \cap \text{space } M \mid A :: - :: \{\text{linorder-topology, second-countable-topology}\} \text{ set. } A \in \text{range greaterThan}\}$ )
by (rule sets-natural-filtration'[OF borel-Ioi])

```

lemma sets-natural-filtration-io:

```

  sets (natural-filtration  $M\ t_0\ X\ t$ ) = sigma-sets (space  $M$ ) ( $\bigcup_{i \in \{t_0..t\}} \{X\ i - ' A \cap \text{space } M \mid A :: - :: \{\text{linorder-topology, second-countable-topology}\} \text{ set. } A \in \text{range lessThan}\}$ )
by (rule sets-natural-filtration'[OF borel-Iio])

```

lemma sets-natural-filtration-ci:

```

  sets (natural-filtration  $M\ t_0\ X\ t$ ) = sigma-sets (space  $M$ ) ( $\bigcup_{i \in \{t_0..t\}} \{X\ i - ' A \cap \text{space } M \mid A :: \text{real set. } A \in \text{range atLeast}\}$ )
by (rule sets-natural-filtration'[OF borel-Ici])

```

lemma (in stochastic-process) subalgebra-natural-filtration:

```

  shows subalgebra  $M$  (natural-filtration  $M\ t_0\ X\ i$ )
  unfolding subalgebra-def using measurable-family-iff-sets by (force simp add:
  natural-filtration-def)

```

sublocale stochastic-process \subseteq filtered-measure-natural-filtration: filtered-measure M natural-filtration $M\ t_0\ X\ t_0$

by (unfold-locales) (intro subalgebra-natural-filtration, simp only: sets-natural-filtration, intro sigma-sets-subseteq, force)

In order to show that the natural filtration constitutes a filtered sigma finite measure, we need to provide a countable exhausting set in the preimage of $X\ t_0$.

lemma (in sigma-finite-measure) sigma-finite-filtered-measure-natural-filtration:

assumes *stochastic-process* $M \ t_0 \ X$
and *exhausting-set*: $\text{countable } A \ (\bigcup A) = \text{space } M \wedge a. a \in A \implies \text{emeasure } M \ a \neq \infty \wedge a. a \in A \implies \exists b \in \text{borel}. a = X \ t_0 - ' b \cap \text{space } M$
shows *sigma-finite-filtered-measure* $M \ (\text{natural-filtration } M \ t_0 \ X) \ t_0$
proof (*unfold-locales*)
interpret *stochastic-process* $M \ t_0 \ X$ **by** (*rule assms*)
have $A \subseteq \text{sets} \ (\text{restr-to-subalg } M \ (\text{natural-filtration } M \ t_0 \ X \ t_0))$ **using** *exhausting-set* **by** (*simp add: sets-restr-to-subalg[OF subalgebra-natural-filtration] sets-natural-filtration*)
fast
moreover **have** $\bigcup A = \text{space} \ (\text{restr-to-subalg } M \ (\text{natural-filtration } M \ t_0 \ X \ t_0))$
unfolding *space-restr-to-subalg* **using** *exhausting-set* **by** *simp*
moreover **have** $\forall a \in A. \text{emeasure} \ (\text{restr-to-subalg } M \ (\text{natural-filtration } M \ t_0 \ X \ t_0)) \ a \neq \infty$ **using** *calculation(1) exhausting-set(3)*
by (*auto simp add: sets-restr-to-subalg[OF subalgebra-natural-filtration] emeasure-restr-to-subalg[OF subalgebra-natural-filtration]*)
ultimately show $\exists A. \text{countable } A \wedge A \subseteq \text{sets} \ (\text{restr-to-subalg } M \ (\text{natural-filtration } M \ t_0 \ X \ t_0)) \wedge \bigcup A = \text{space} \ (\text{restr-to-subalg } M \ (\text{natural-filtration } M \ t_0 \ X \ t_0)) \wedge (\forall a \in A. \text{emeasure} \ (\text{restr-to-subalg } M \ (\text{natural-filtration } M \ t_0 \ X \ t_0)) \ a \neq \infty)$ **using** *exhausting-set* **by** *blast*
show $\bigwedge i \ j. \llbracket t_0 \leq i; i \leq j \rrbracket \implies \text{sets} \ (\text{natural-filtration } M \ t_0 \ X \ i) \subseteq \text{sets} \ (\text{natural-filtration } M \ t_0 \ X \ j)$ **using** *filtered-measure-natural-filtration.subalgebra-F* **by** (*simp add: subalgebra-def*)
qed (*auto intro: stochastic-process.subalgebra-natural-filtration assms(1)*)

lemma (*in finite-measure*) *finite-filtered-measure-natural-filtration*:

assumes *stochastic-process* $M \ t_0 \ X$
shows *finite-filtered-measure* $M \ (\text{natural-filtration } M \ t_0 \ X) \ t_0$
proof
interpret *stochastic-process* $M \ t_0 \ X$ **by** (*rule assms*)
show $t_0 \leq i \implies \text{subalgebra } M \ (\text{natural-filtration } M \ t_0 \ X \ i)$ **for** i **using** *subalgebra-natural-filtration* **by** *blast*
show $\llbracket t_0 \leq i; i \leq j \rrbracket \implies \text{sets} \ (\text{natural-filtration } M \ t_0 \ X \ i) \subseteq \text{sets} \ (\text{natural-filtration } M \ t_0 \ X \ j)$ **for** $i \ j$ **using** *filtered-measure-natural-filtration.subalgebra-F* **by** (*simp add: subalgebra-def*)
qed

10.2 Adapted Process

We call a collection a stochastic process X adapted if $X \ i$ is $F \ i$ -borel-measurable for all indices i .

locale *adapted-process* = *filtered-measure* $M \ F \ t_0$ **for** $M \ F \ t_0$ **and** $X :: - \Rightarrow - \Rightarrow - :: \{\text{second-countable-topology, banach}\} +$
assumes *adapted[measurable]*: $\bigwedge i. t_0 \leq i \implies X \ i \in \text{borel-measurable} \ (F \ i)$
begin

lemma *adaptedE[elim]*:

assumes $\llbracket \bigwedge i. t_0 \leq j \implies j \leq i \implies X \ j \in \text{borel-measurable} \ (F \ i) \rrbracket \implies P$
shows P
using *assms* **using** *adapted* **by** (*metis dual-order.trans borel-measurable-subalgebra*)

sets-F-mono space-F)

lemma *adaptedD*:

assumes $t_0 \leq j \leq i$

shows $X j \in \text{borel-measurable } (F i)$ **using** *assms adaptedE* **by** *meson*

end

locale *nat-adapted-process* = *adapted-process* $M F 0 :: \text{nat } X$ **for** $M F X$

locale *real-adapted-process* = *adapted-process* $M F 0 :: \text{real } X$ **for** $M F X$

sublocale *nat-adapted-process* \subseteq *nat-filtered-measure* ..

sublocale *real-adapted-process* \subseteq *real-filtered-measure* ..

lemma (*in filtered-measure*) *adapted-process-const-fun*:

assumes $f \in \text{borel-measurable } (F t_0)$

shows *adapted-process* $M F t_0 (\lambda \cdot. f)$

using *measurable-from-subalg subalgebra-F assms* **by** (*unfold-locales*) *blast*

lemma (*in filtered-measure*) *adapted-process-const*:

shows *adapted-process* $M F t_0 (\lambda i \cdot. c i)$ **by** (*unfold-locales*) *simp*

context *adapted-process*

begin

lemma *compose-adapted*:

assumes $\bigwedge i. t_0 \leq i \implies f i \in \text{borel-measurable borel}$

shows *adapted-process* $M F t_0 (\lambda i \xi. (f i) (X i \xi))$

by (*unfold-locales*) (*intro measurable-compose[OF adapted assms]*)

lemma *norm-adapted*: *adapted-process* $M F t_0 (\lambda i \xi. \text{norm } (X i \xi))$ **by** (*fastforce intro: compose-adapted*)

lemma *scaleR-right-adapted*:

assumes *adapted-process* $M F t_0 R$

shows *adapted-process* $M F t_0 (\lambda i \xi. (R i \xi) *_R (X i \xi))$

using *adapted-process.adapted[OF assms]* *adapted* **by** (*unfold-locales*) *simp*

lemma *scaleR-right-const-fun-adapted*:

assumes $f \in \text{borel-measurable } (F t_0)$

shows *adapted-process* $M F t_0 (\lambda i \xi. f \xi *_R (X i \xi))$

using *assms* **by** (*fast intro: scaleR-right-adapted adapted-process-const-fun*)

lemma *scaleR-right-const-adapted*: *adapted-process* $M F t_0 (\lambda i \xi. c i *_R (X i \xi))$ **by** (*unfold-locales*) *simp*

lemma *add-adapted*:

assumes *adapted-process* $M F t_0 Y$

shows *adapted-process* $M F t_0 (\lambda i \xi. X i \xi + Y i \xi)$

using *adapted-process.adapted*[*OF assms*] **adapted by** (*unfold-locales*) *simp*

lemma *diff-adapted*:
assumes *adapted-process* $M F t_0 Y$
shows *adapted-process* $M F t_0 (\lambda i \xi. X i \xi - Y i \xi)$
using *adapted-process.adapted*[*OF assms*] **adapted by** (*unfold-locales*) *simp*

lemma *uminus-adapted*: *adapted-process* $M F t_0 (-X)$ **using** *scaleR-right-const-adapted*[*of* $\lambda -. -1$] **by** (*simp add: fun-Compl-def*)

lemma *partial-sum-adapted*: *adapted-process* $M F t_0 (\lambda n \xi. \sum_{i \in \{t_0..n\}} X i \xi)$
proof (*unfold-locales*)
fix $i :: 'b$
have $X j \in \text{borel-measurable } (F i)$ **if** $t_0 \leq j$ **for** $j \leq i$ **for** j **using** *that adaptedE* **by** *meson*
thus $(\lambda \xi. \sum_{i \in \{t_0..i\}} X i \xi) \in \text{borel-measurable } (F i)$ **by** *simp*
qed

lemma *partial-sum'-adapted*: *adapted-process* $M F t_0 (\lambda n \xi. \sum_{i \in \{t_0..<n\}} X i \xi)$
proof (*unfold-locales*)
fix $i :: 'b$
have $X j \in \text{borel-measurable } (F i)$ **if** $t_0 \leq j$ **for** $j < i$ **for** j **using** *that adaptedE* **by** *fastforce*
thus $(\lambda \xi. \sum_{i \in \{t_0..<i\}} X i \xi) \in \text{borel-measurable } (F i)$ **by** *simp*
qed

end

lemma (*in nat-adapted-process*) *partial-sum-Suc-adapted*: *nat-adapted-process* $M F (\lambda n \xi. \sum_{i < n} X (Suc i) \xi)$
proof (*unfold-locales*)
fix i
have $X j \in \text{borel-measurable } (F i)$ **if** $j \leq i$ **for** j **using** *that adaptedD* **by** *blast*
thus $(\lambda \xi. \sum_{i < i} X (Suc i) \xi) \in \text{borel-measurable } (F i)$ **by** *auto*
qed

lemma (*in filtered-measure*) *adapted-process-sum*:
assumes $\bigwedge i. i \in I \implies \text{adapted-process } M F t_0 (X i)$
shows *adapted-process* $M F t_0 (\lambda k \xi. \sum_{i \in I} X i k \xi)$
proof –
{
fix $i k$ **assume** $i \in I$ **and** *asm*: $t_0 \leq k$
then interpret *adapted-process* $M F t_0 X i$ **using** *assms* **by** *simp*
have $X i k \in \text{borel-measurable } M X i k \in \text{borel-measurable } (F k)$ **using** *measurable-from-subalg subalgebra adapted asm* **by** (*blast, simp*)
}
thus *?thesis* **by** (*unfold-locales*) *simp*
qed

An adapted process is necessarily a stochastic process.

sublocale *adapted-process* \subseteq *stochastic-process* **using** *measurable-from-subalg sub-algebra adapted* **by** (*unfold-locales*) *blast*

sublocale *nat-adapted-process* \subseteq *nat-stochastic-process* **..**

sublocale *real-adapted-process* \subseteq *real-stochastic-process* **..**

A stochastic process is always adapted to the natural filtration it generates.

sublocale *stochastic-process* \subseteq *adapted-natural: adapted-process M natural-filtration M t₀ X t₀ X* **by** (*unfold-locales*) (*auto simp add: natural-filtration-def intro: random-variable measurable-family-vimage-algebra*)

10.3 Progressively Measurable Process

locale *progressive-process* = *filtered-measure M F t₀ for M F t₀ and X :: - \Rightarrow - \Rightarrow - :: {second-countable-topology, banach} +*

assumes *progressive[measurable]: $\bigwedge t. t_0 \leq t \implies (\lambda(i, x). X i x) \in \text{borel-measurable}(\text{restrict-space borel } \{t_0..t\} \otimes_M F t)$*

begin

lemma *progressiveD:*

assumes *S \in borel*

shows *($\lambda(j, \xi). X j \xi$) - ' S \cap ($\{t_0..i\} \times \text{space } M$) \in (restrict-space borel $\{t_0..i\} \otimes_M F i$)*

using *measurable-sets[OF progressive, OF - assms, of i]*

by (*cases t₀ \leq i*) (*auto simp add: space-restrict-space sets-pair-measure space-pair-measure*)

end

locale *nat-progressive-process* = *progressive-process M F 0 :: nat X for M F X*

locale *real-progressive-process* = *progressive-process M F 0 :: real X for M F X*

lemma (*in filtered-measure*) *progressive-process-const-fun:*

assumes *f \in borel-measurable (F t₀)*

shows *progressive-process M F t₀ ($\lambda \cdot. f$)*

proof (*unfold-locales*)

fix *i* **assume** *asm: t₀ \leq i*

have *f \in borel-measurable (F i) using borel-measurable-mono[OF order.refl asm] assms by blast*

thus *case-prod ($\lambda \cdot. f$) \in borel-measurable (restrict-space borel $\{t_0..i\} \otimes_M F i$)*

using *measurable-compose[OF measurable-snd] by simp*

qed

lemma (*in filtered-measure*) *progressive-process-const:*

assumes *c \in borel-measurable borel*

shows *progressive-process M F t₀ ($\lambda i \cdot. c i$)*

using *assms by (unfold-locales) (auto simp add: measurable-split-conv intro!: measurable-compose[OF measurable-fst] measurable-restrict-space1)*

context *progressive-process*
begin

lemma *compose-progressive:*

assumes *case-prod* $f \in \text{borel-measurable borel}$
shows *progressive-process* $M F t_0 (\lambda i \xi. (f i) (X i \xi))$
proof
fix i **assume** *asm*: $t_0 \leq i$
have $(\lambda(j, \xi). (j, X j \xi)) \in (\text{restrict-space borel } \{t_0..i\} \otimes_M F i) \rightarrow_M \text{borel} \otimes_M \text{borel}$
using *progressive*[*OF asm*] *measurable-fst'*[*OF measurable-restrict-space1*, *OF measurable-id*]
by (*auto simp add: measurable-pair-iff measurable-split-conv*)
moreover have $(\lambda(j, \xi). f j (X j \xi)) = \text{case-prod } f \circ ((\lambda(j, y). (j, y)) \circ (\lambda(j, \xi). (j, X j \xi)))$ **by** *fastforce*
ultimately show $(\lambda(j, \xi). (f j) (X j \xi)) \in \text{borel-measurable } (\text{restrict-space borel } \{t_0..i\} \otimes_M F i)$ **using** *assms* **by** (*simp add: borel-prod*)
qed

lemma *norm-progressive: progressive-process* $M F t_0 (\lambda i \xi. \text{norm } (X i \xi))$ **using** *measurable-compose*[*OF progressive borel-measurable-norm*] **by** (*unfold-locales*) *simp*

lemma *scaleR-right-progressive:*

assumes *progressive-process* $M F t_0 R$
shows *progressive-process* $M F t_0 (\lambda i \xi. (R i \xi) *_R (X i \xi))$
using *progressive-process.progressive*[*OF assms*] **by** (*unfold-locales*) (*simp add: progressive assms*)

lemma *scaleR-right-const-fun-progressive:*

assumes $f \in \text{borel-measurable } (F t_0)$
shows *progressive-process* $M F t_0 (\lambda i \xi. f \xi *_R (X i \xi))$
using *assms* **by** (*fast intro: scaleR-right-progressive progressive-process-const-fun*)

lemma *scaleR-right-const-progressive:*

assumes $c \in \text{borel-measurable borel}$
shows *progressive-process* $M F t_0 (\lambda i \xi. c i *_R (X i \xi))$
using *assms* **by** (*fastforce intro: scaleR-right-progressive progressive-process-const*)

lemma *add-progressive:*

assumes *progressive-process* $M F t_0 Y$
shows *progressive-process* $M F t_0 (\lambda i \xi. X i \xi + Y i \xi)$
using *progressive-process.progressive*[*OF assms*] **by** (*unfold-locales*) (*simp add: progressive assms*)

lemma *diff-progressive:*

assumes *progressive-process* $M F t_0 Y$
shows *progressive-process* $M F t_0 (\lambda i \xi. X i \xi - Y i \xi)$
using *progressive-process.progressive*[*OF assms*] **by** (*unfold-locales*) (*simp add:*

progressive assms)

lemma *uminus-progressive*: *progressive-process* $M \ F \ t_0 \ (-X)$ **using** *scaleR-right-const-progressive*[*of* $\lambda \cdot, -1$] **by** (*simp add: fun-Compl-def*)

end

A progressively measurable process is also adapted.

sublocale *progressive-process* \subseteq *adapted-process* **using** *measurable-compose-rev*[*OF* *progressive measurable-Pair1*]

unfolding *prod.case space-restrict-space*

by *unfold-locales simp*

sublocale *nat-progressive-process* \subseteq *nat-adapted-process* ..

sublocale *real-progressive-process* \subseteq *real-adapted-process* ..

In the discrete setting, adaptedness is equivalent to progressive measurability.

sublocale *nat-adapted-process* \subseteq *nat-progressive-process*

proof (*unfold-locales, intro borel-measurableI*)

fix $S :: 'b \text{ set}$ **and** $i :: \text{nat}$ **assume** *open-S: open S*

{

fix j **assume** *asm: $j \leq i$*

hence $X \ j - ' S \cap \text{space } M \in F \ i$ **using** *adaptedD[of j, THEN measurable-sets]*
space-F open-S **by** *fastforce*

moreover have *case-prod* $X - ' S \cap \{j\} \times \text{space } M = \{j\} \times (X \ j - ' S \cap \text{space } M)$ **for** j **by** *fast*

moreover have $\{j :: \text{nat}\} \in \text{restrict-space borel } \{0..i\}$ **using** *asm* **by** (*simp add: sets-restrict-space-iff*)

ultimately have *case-prod* $X - ' S \cap \{j\} \times \text{space } M \in \text{restrict-space borel } \{0..i\} \otimes_M F \ i$ **by** *simp*

}

hence $(\lambda j. (\lambda(x, y). X \ x \ y) - ' S \cap \{j\} \times \text{space } M) \ ' \{..i\} \subseteq \text{restrict-space borel } \{0..i\} \otimes_M F \ i$ **by** *blast*

moreover have *case-prod* $X - ' S \cap \text{space } (\text{restrict-space borel } \{0..i\} \otimes_M F \ i) = (\bigcup_{j \leq i} \text{case-prod } X - ' S \cap \{j\} \times \text{space } M)$ **unfolding** *space-pair-measure space-restrict-space space-F* **by** *force*

ultimately show *case-prod* $X - ' S \cap \text{space } (\text{restrict-space borel } \{0..i\} \otimes_M F \ i) \in \text{restrict-space borel } \{0..i\} \otimes_M F \ i$ **by** (*metis sets.countable-UN*)

qed

10.4 Predictable Process

We introduce the constant Σ_P to denote the predictable sigma algebra.

context *linearly-filtered-measure*

begin

definition $\Sigma_P :: ('b \times 'a) \text{ measure where predictable-sigma: } \Sigma_P \equiv \text{sigma } (\{t_0..\} \times \text{space } M) (\{\{s<..t\} \times A \mid A \text{ s t. } A \in F \text{ s } \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in F \text{ } t_0\})$

lemma *space-predictable-sigma[simp]: space $\Sigma_P = (\{t_0..\} \times \text{space } M)$ unfolding predictable-sigma space-measure-of-conv by blast*

lemma *sets-predictable-sigma: sets $\Sigma_P = \text{sigma-sets } (\{t_0..\} \times \text{space } M) (\{\{s<..t\} \times A \mid A \text{ s t. } A \in F \text{ s } \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in F \text{ } t_0\})$ unfolding predictable-sigma using space-F sets.sets-into-space by (subst sets-measure-of) fastforce+*

lemma *measurable-predictable-sigma-snd:*
assumes *countable \mathcal{I} $\mathcal{I} \subseteq \{\{s<..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\} \{t_0<..\} \subseteq (\bigcup \mathcal{I})$*
shows *snd $\in \Sigma_P \rightarrow_M F \text{ } t_0$*
proof (intro measurableI, force)
fix *S :: 'a set assume asm: $S \in F \text{ } t_0$*
have *countable: countable $((\lambda I. I \times S) \text{ ' } \mathcal{I})$ using assms(1) by blast*
have *$(\lambda I. I \times S) \text{ ' } \mathcal{I} \subseteq \{\{s<..t\} \times A \mid A \text{ s t. } A \in F \text{ s } \wedge t_0 \leq s \wedge s < t\}$ using sets-F-mono[OF order-refl, THEN subsetD, OF - asm] assms(2) by blast*
hence *$(\bigcup I \in \mathcal{I}. I \times S) \cup \{t_0\} \times S \in \Sigma_P$ unfolding sets-predictable-sigma using asm by (intro sigma-sets-Un[OF sigma-sets-UNION[OF countable] sigma-sets.Basic] sigma-sets.Basic) blast+*
moreover have *snd - ' $S \cap \text{space } \Sigma_P = \{t_0..\} \times S$ using sets.sets-into-space[OF asm] by fastforce*
moreover have *$\{t_0\} \cup \{t_0<..\} = \{t_0..\}$ by auto*
moreover have *$(\bigcup I \in \mathcal{I}. I \times S) \cup \{t_0\} \times S = \{t_0..\} \times S$ using assms(2,3) calculation(3) by fastforce*
ultimately show *snd - ' $S \cap \text{space } \Sigma_P \in \Sigma_P$ by argo*
qed

lemma *measurable-predictable-sigma-fst:*
assumes *countable \mathcal{I} $\mathcal{I} \subseteq \{\{s<..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\} \{t_0<..\} \subseteq (\bigcup \mathcal{I})$*
shows *fst $\in \Sigma_P \rightarrow_M \text{borel}$*
proof -
have *$A \times \text{space } M \in \text{sets } \Sigma_P$ if $A \in \text{sigma-sets } \{t_0..\} \{\{s<..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\}$ for A unfolding sets-predictable-sigma using that*
proof (induction rule: sigma-sets.induct)
case (Basic a)
thus *?case using space-F sets.top by blast*
next
case (Compl a)
have *$(\{t_0..\} - a) \times \text{space } M = \{t_0..\} \times \text{space } M - a \times \text{space } M$ by blast*
then show *?case using Compl(2)[THEN sigma-sets.Compl] by presburger*
next
case (Union a)
have *$\bigcup (\text{range } a) \times \text{space } M = \bigcup (\text{range } (\lambda i. a \text{ } i \times \text{space } M))$ by blast*
then show *?case using Union(2)[THEN sigma-sets.Union] by presburger*
qed (auto)

```

moreover have restrict-space borel  $\{t_0..\} = \text{sigma } \{t_0..\} \{ \{s<..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t \}$ 
proof –
  have sigma-sets  $\{t_0..\} ((\cap) \{t_0..\} \text{ ‘sigma-sets UNIV (range greaterThan)}) =$ 
sigma-sets  $\{t_0..\} \{ \{s<..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t \}$ 
proof (intro sigma-sets-eqI ; clarify)
  fix A :: 'b set assume asm:  $A \in \text{sigma-sets UNIV (range greaterThan)}$ 
  thus  $\{t_0..\} \cap A \in \text{sigma-sets } \{t_0..\} \{ \{s<..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t \}$ 
proof (induction rule: sigma-sets.induct)
  case (Basic a)
  then obtain s where  $s: a = \{s<..\}$  by blast
  show ?case
  proof (cases  $t_0 \leq s$ )
  case True
  hence *:  $\{t_0..\} \cap a = (\bigcup i \in \mathcal{I}. \{s<..\} \cap i)$  using s assms(3) by force
  have  $((\cap) \{s<..\} \text{ ‘}\mathcal{I}) \subseteq \text{sigma-sets } \{t_0..\} \{ \{s<..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t \}$ 
proof (clarify)
  fix A assume  $A \in \mathcal{I}$ 
  then obtain s' t' where  $A: A = \{s'<..t'\} \ t_0 \leq s' \ s' < t'$  using assms(2)
by blast
  hence  $\{s<..\} \cap A = \{ \max s \ s'<..t' \}$  by fastforce
  moreover have  $t_0 \leq \max s \ s'$  using A True by linarith
  moreover have  $\max s \ s' < t'$  if  $s < t'$  using A that by linarith
  moreover have  $\{s<..\} \cap A = \{ \}$  if  $\neg s < t'$  using A that by force
  ultimately show  $\{s<..\} \cap A \in \text{sigma-sets } \{t_0..\} \{ \{s<..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t \}$  by (cases  $s < t'$ ) (blast, simp add: sigma-sets.Empty)
  qed
  thus ?thesis unfolding * using assms(1) by (intro sigma-sets-UNION)
auto
next
case False
  hence  $\{t_0..\} \cap a = \{t_0..\}$  using s by force
  thus ?thesis using sigma-sets-top by auto
qed
next
case (Compl a)
  have  $\{t_0..\} \cap (\text{UNIV} - a) = \{t_0..\} - (\{t_0..\} \cap a)$  by blast
  then show ?case using Compl(2)[THEN sigma-sets.Compl] by presburger
next
case (Union a)
  have  $\{t_0..\} \cap \bigcup (\text{range } a) = \bigcup (\text{range } (\lambda i. \{t_0..\} \cap a \ i))$  by blast
  then show ?case using Union(2)[THEN sigma-sets.Union] by presburger
qed (simp add: sigma-sets.Empty)
next
fix s t assume asm:  $t_0 \leq s \ s < t$ 
  hence *:  $\{s<..t\} = \{s<..\} \cap (\{t_0..\} - \{t<..\})$  by force
  have  $\{s<..\} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\} \text{ ‘sigma-sets UNIV (range greaterThan)})$ 
using asm by (intro sigma-sets.Basic) auto
  moreover have  $\{t_0..\} - \{t<..\} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\} \text{ ‘sigma-sets$ 

```

```

UNIV (range greaterThan)) using asm by (intro sigma-sets.Compl sigma-sets.Basic)
auto
  ultimately show  $\{s <.. t\} \in \text{sigma-sets } \{t_{0..}\} ((\cap) \{t_{0..}\} \text{ ' sigma-sets UNIV$ 
(range greaterThan)) unfolding * Int-range-binary[of  $\{s <.. \}$ ] by (intro sigma-sets-Inter[OF
- binary-in-sigma-sets]) auto
  qed
  thus ?thesis unfolding borel-Ioi restrict-space-def emeasure-sigma by (force
intro: sigma-eqI)
  qed
  ultimately have restrict-space borel  $\{t_{0..}\} \otimes_M \text{sigma (space } M) \{ \} \subseteq \text{sets } \Sigma_P$ 

    unfolding sets-pair-measure space-restrict-space space-measure-of-conv
    using space-predictable-sigma sets.sigma-algebra-axioms[of  $\Sigma_P$ ]
    by (intro sigma-algebra.sigma-sets-subset) (auto simp add: sigma-sets-empty-eq
sets-measure-of-conv)
    moreover have space (restrict-space borel  $\{t_{0..}\} \otimes_M \text{sigma (space } M) \{ \}) =$ 
space  $\Sigma_P$  by (simp add: space-pair-measure)
    moreover have  $\text{fst} \in \text{restrict-space borel } \{t_{0..}\} \otimes_M \text{sigma (space } M) \{ \} \rightarrow_M$ 
borel by (fastforce intro: measurable-fst''[OF measurable-restrict-space1, of  $\lambda x. x$ ])

    ultimately show ?thesis by (meson borel-measurable-subalgebra)
  qed
end

locale predictable-process = linearly-filtered-measure M F t0 for M F t0 and X ::
-  $\Rightarrow$  -  $\Rightarrow$  - :: {second-countable-topology, banach} +
  assumes predictable:  $(\lambda(t, x). X t x) \in \text{borel-measurable } \Sigma_P$ 
begin

lemmas predictableD = measurable-sets[OF predictable, unfolded space-predictable-sigma]

end

locale nat-predictable-process = predictable-process M F 0 :: nat X for M F X
locale real-predictable-process = predictable-process M F 0 :: real X for M F X

lemma (in nat-filtered-measure) measurable-predictable-sigma-snd':
  shows  $\text{snd} \in \Sigma_P \rightarrow_M F 0$ 
  by (intro measurable-predictable-sigma-snd[of range  $(\lambda x. \{ \text{Suc } x \})$ ]) (force | simp
add: greaterThan-0)+

lemma (in nat-filtered-measure) measurable-predictable-sigma-fst':
  shows  $\text{fst} \in \Sigma_P \rightarrow_M \text{borel}$ 
  by (intro measurable-predictable-sigma-fst[of range  $(\lambda x. \{ \text{Suc } x \})$ ]) (force | simp
add: greaterThan-0)+

```

lemma (in *real-filtered-measure*) *measurable-predictable-sigma-snd'*:
 shows $snd \in \Sigma_P \rightarrow_M F\ 0$
 using *real-arch-simple* **by** (intro *measurable-predictable-sigma-snd* [of range $(\lambda x::nat. \{0 < ..real\ (Suc\ x)\})$]) (fastforce intro: *add-increasing*) +

lemma (in *real-filtered-measure*) *measurable-predictable-sigma-fst'*:
 shows $fst \in \Sigma_P \rightarrow_M borel$
 using *real-arch-simple* **by** (intro *measurable-predictable-sigma-fst* [of range $(\lambda x::nat. \{0 < ..real\ (Suc\ x)\})$]) (fastforce intro: *add-increasing*) +

lemma (in *linearly-filtered-measure*) *predictable-process-const-fun*:
 assumes $snd \in \Sigma_P \rightarrow_M F\ t_0$ $f \in borel\text{-}measurable\ (F\ t_0)$
 shows *predictable-process* $M\ F\ t_0\ (\lambda -. f)$
 using *measurable-compose-rev* [OF *assms*(2)] *assms*(1) **by** (*unfold-locales*) (*auto simp add: measurable-split-conv*)

lemma (in *nat-filtered-measure*) *predictable-process-const-fun'* [intro]:
 assumes $f \in borel\text{-}measurable\ (F\ 0)$
 shows *nat-predictable-process* $M\ F\ (\lambda -. f)$
 using *assms* **by** (intro *predictable-process-const-fun* [OF *measurable-predictable-sigma-snd'*, *THEN nat-predictable-process.intro*])

lemma (in *real-filtered-measure*) *predictable-process-const-fun'* [intro]:
 assumes $f \in borel\text{-}measurable\ (F\ 0)$
 shows *real-predictable-process* $M\ F\ (\lambda -. f)$
 using *assms* **by** (intro *predictable-process-const-fun* [OF *measurable-predictable-sigma-snd'*, *THEN real-predictable-process.intro*])

lemma (in *linearly-filtered-measure*) *predictable-process-const*:
 assumes $fst \in borel\text{-}measurable\ \Sigma_P$ $c \in borel\text{-}measurable\ borel$
 shows *predictable-process* $M\ F\ t_0\ (\lambda i -. c\ i)$
 using *assms* **by** (*unfold-locales*) (*simp add: measurable-split-conv*)

lemma (in *linearly-filtered-measure*) *predictable-process-const-const* [intro]:
 shows *predictable-process* $M\ F\ t_0\ (\lambda -. c)$
by (*unfold-locales*) *simp*

lemma (in *nat-filtered-measure*) *predictable-process-const'* [intro]:
 assumes $c \in borel\text{-}measurable\ borel$
 shows *nat-predictable-process* $M\ F\ (\lambda i -. c\ i)$
 using *assms* **by** (intro *predictable-process-const* [OF *measurable-predictable-sigma-fst'*, *THEN nat-predictable-process.intro*])

lemma (in *real-filtered-measure*) *predictable-process-const'* [intro]:
 assumes $c \in borel\text{-}measurable\ borel$
 shows *real-predictable-process* $M\ F\ (\lambda i -. c\ i)$
 using *assms* **by** (intro *predictable-process-const* [OF *measurable-predictable-sigma-fst'*,

THEN real-predictable-process.intro])

context *predictable-process*
begin

lemma *compose-predictable:*

assumes *fst* \in *borel-measurable* Σ_P *case-prod* *f* \in *borel-measurable borel*
shows *predictable-process* *M F t₀* $(\lambda i \xi. (f i) (X i \xi))$

proof

have $(\lambda(i, \xi). (i, X i \xi)) \in \Sigma_P \rightarrow_M \text{borel} \otimes_M \text{borel}$ **using** *predictable assms(1)*
by (*auto simp add: measurable-pair-iff measurable-split-conv*)

moreover have $(\lambda(i, \xi). f i (X i \xi)) = \text{case-prod } f \circ (\lambda(i, \xi). (i, X i \xi))$ **by**
fastforce

ultimately show $(\lambda(i, \xi). f i (X i \xi)) \in \text{borel-measurable } \Sigma_P$ **unfolding** *borel-prod*
using *assms* **by** *simp*

qed

lemma *norm-predictable:* *predictable-process* *M F t₀* $(\lambda i \xi. \text{norm } (X i \xi))$ **using**
measurable-compose[OF predictable borel-measurable-norm]

by (*unfold-locales*) (*simp add: prod.case-distrib*)

lemma *scaleR-right-predictable:*

assumes *predictable-process* *M F t₀* *R*
shows *predictable-process* *M F t₀* $(\lambda i \xi. (R i \xi) *_R (X i \xi))$

using *predictable predictable-process.predictable[OF assms]* **by** (*unfold-locales*)
(auto simp add: measurable-split-conv)

lemma *scaleR-right-const-fun-predictable:*

assumes *snd* $\in \Sigma_P \rightarrow_M F t_0$ *f* \in *borel-measurable (F t₀)*

shows *predictable-process* *M F t₀* $(\lambda i \xi. f \xi *_R (X i \xi))$

using *assms* **by** (*fast intro: scaleR-right-predictable predictable-process-const-fun*)

lemma *scaleR-right-const-predictable:*

assumes *fst* \in *borel-measurable* Σ_P *c* \in *borel-measurable borel*

shows *predictable-process* *M F t₀* $(\lambda i \xi. c i *_R (X i \xi))$

using *assms* **by** (*fastforce intro: scaleR-right-predictable predictable-process-const*)

lemma *scaleR-right-const'-predictable:* *predictable-process* *M F t₀* $(\lambda i \xi. c *_R (X i \xi))$

by (*fastforce intro: scaleR-right-predictable predictable-process-const-const*)

lemma *add-predictable:*

assumes *predictable-process* *M F t₀* *Y*

shows *predictable-process* *M F t₀* $(\lambda i \xi. X i \xi + Y i \xi)$

using *predictable predictable-process.predictable[OF assms]* **by** (*unfold-locales*)
(auto simp add: measurable-split-conv)

lemma *diff-predictable:*

assumes *predictable-process* *M F t₀* *Y*

shows *predictable-process* $M F t_0 (\lambda i \xi. X i \xi - Y i \xi)$
using *predictable* *predictable-process.predictable*[*OF assms*] **by** (*unfold-locales*)
(auto simp add: measurable-split-conv)

lemma *uminus-predictable*: *predictable-process* $M F t_0 (-X)$ **using** *scaleR-right-const'-predictable*[*of -1*] **by** (*simp add: fun-Compl-def*)

end

Every predictable process is also progressively measurable.

sublocale *predictable-process* \subseteq *progressive-process*

proof (*unfold-locales*)

fix $i :: 'b$ **assume** $asm: t_0 \leq i$
 $\{$
fix $S :: ('b \times 'a)$ *set* **assume** $S \in \{\{s <..t\} \times A \mid A s t. A \in F s \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in F t_0\}$
hence $(\lambda x. x) - ' S \cap (\{t_0..i\} \times \text{space } M) \in \text{restrict-space borel } \{t_0..i\} \otimes_M F$
 i

proof

assume $S \in \{\{s <..t\} \times A \mid A s t. A \in F s \wedge t_0 \leq s \wedge s < t\}$
then obtain $s t A$ **where** $S\text{-is}: S = \{s <..t\} \times A$ $t_0 \leq s$ $s < t$ $A \in F s$ **by**
blast

hence $(\lambda x. x) - ' S \cap (\{t_0..i\} \times \text{space } M) = \{s <.. \min i t\} \times A$ **using**
sets.sets-into-space[*OF S-is(4)*] **by** *auto*

then show *?thesis* **using** $S\text{-is}$ *sets-F-mono*[*of s i*] **by** (*cases s ≤ i*) (*fastforce simp add: sets-restrict-space-iff*) +

next

assume $S \in \{\{t_0\} \times A \mid A. A \in F t_0\}$
then obtain A **where** $S\text{-is}: S = \{t_0\} \times A$ $A \in F t_0$ **by** *blast*
hence $(\lambda x. x) - ' S \cap (\{t_0..i\} \times \text{space } M) = \{t_0\} \times A$ **using** asm *sets.sets-into-space*[*OF S-is(2)*] **by** *auto*

thus *?thesis* **using** $S\text{-is}(2)$ *sets-F-mono*[*OF order-refl asm*] asm **by** (*fastforce simp add: sets-restrict-space-iff*)

qed

hence $(\lambda x. x) - ' S \cap \text{space } (\text{restrict-space borel } \{t_0..i\} \otimes_M F i) \in \text{restrict-space borel } \{t_0..i\} \otimes_M F i$ **by** (*simp add: space-pair-measure space-F*[*OF asm*])
 $\}$

moreover have $\{\{s <..t\} \times A \mid A s t. A \in \text{sets } (F s) \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in \text{sets } (F t_0)\} \subseteq \text{Pow } (\{t_0..i\} \times \text{space } M)$ **using** *sets.sets-into-space* **by**
force

ultimately have $(\lambda x. x) \in \text{restrict-space borel } \{t_0..i\} \otimes_M F i \rightarrow_M \Sigma_P$ **using** *space-F*[*OF asm*] **by** (*intro measurable-sigma-sets*[*OF sets-predictable-sigma*])
(fast, force simp add: space-pair-measure)

thus case-prod $X \in \text{borel-measurable } (\text{restrict-space borel } \{t_0..i\} \otimes_M F i)$ **using**
predictable **by** *simp*

qed

sublocale *nat-predictable-process* \subseteq *nat-progressive-process* ..

sublocale *real-predictable-process* \subseteq *real-progressive-process* ..

The following lemma characterizes predictability in a discrete-time setting.

lemma (in *nat-filtered-measure*) *sets-in-filtration*:

assumes $(\bigcup i. \{i\} \times A \ i) \in \Sigma_P$

shows $A \ (Suc \ i) \in F \ i \ A \ 0 \in F \ 0$

using *assms unfolding sets-predictable-sigma*

proof (*induction* $(\bigcup i. \{i\} \times A \ i)$ *arbitrary: A*)

case *Basic*

{

assume $\exists S. (\bigcup i. \{i\} \times A \ i) = \{0\} \times S$

then obtain S **where** $S: (\bigcup i. \{i\} \times A \ i) = \{bot\} \times S$ **unfolding** *bot-nat-def*

by *blast*

hence $S \in F \ bot$ **using** *Basic* **by** (*fastforce simp add: times-eq-iff bot-nat-def*)

moreover have $A \ i = \{\}$ **if** $i \neq bot$ **for** i **using** *that S* **by** *blast*

moreover have $A \ bot = S$ **using** S **by** *blast*

ultimately have $A \ (Suc \ i) \in F \ i \ A \ 0 \in F \ 0$ **for** i **unfolding** *bot-nat-def* **by** (*auto simp add: bot-nat-def*)

}

note $*$ = *this*

{

assume $\nexists S. (\bigcup i. \{i\} \times A \ i) = \{0\} \times S$

then obtain $s \ t \ B$ **where** $B: (\bigcup i. \{i\} \times A \ i) = \{s<..t\} \times B$ $B \in sets \ (F \ s)$

$s < t$ **using** *Basic* **by** *auto*

hence $A \ i = B$ **if** $i \in \{s<..t\}$ **for** i **using** *that* **by** *fast*

moreover have $A \ i = \{\}$ **if** $i \notin \{s<..t\}$ **for** i **using** B **that** **by** *fastforce*

ultimately have $A \ (Suc \ i) \in F \ i \ A \ 0 \in F \ 0$ **for** i **unfolding** *bot-nat-def* **using** B *sets-F-mono* **by** (*auto simp add: bot-nat-def*) (*metis less-Suc-eq-le sets.empty-sets subset-eq*)

}

note $**$ = *this*

show $A \ (Suc \ i) \in sets \ (F \ i) \ A \ 0 \in sets \ (F \ 0)$ **using** $*(1)[of \ i] \ *(2) \ **(1)[of \ i]$

$**(2)$ **by** *blast+*

next

case *Empty*

{

case *1*

then show *?case* **using** *Empty* **by** *simp*

next

case *2*

then show *?case* **using** *Empty* **by** *simp*

}

next

case (*Compl a*)

have $a\text{-in}: a \subseteq \{0..\} \times space \ M$ **using** *Compl(1) sets.sets-into-space sets-predictable-sigma space-predictable-sigma* **by** *metis*

hence $A\text{-in}: A \ i \subseteq space \ M$ **for** i **using** *Compl(4)* **by** *blast*

have $a: a = \{0..\} \times space \ M - (\bigcup i. \{i\} \times A \ i)$ **using** $a\text{-in}$ *Compl(4)* **by** *blast*

also have $... = - (\bigcap j. - (\{j\} \times (space \ M - A \ j)))$ **by** *blast*

also have $... = (\bigcup j. \{j\} \times (space \ M - A \ j))$ **by** *blast*

finally have $*$: $(space \ M - A \ (Suc \ i)) \in F \ i \ (space \ M - A \ 0) \in F \ 0$ **using**

```

Compl(2,3) by auto
{
  case 1
    then show ?case using * A-in by (metis bot-nat-0.extremum double-diff
sets.Diff sets.top sets-F-mono sets-le-imp-space-le space-F)
  next
    case 2
      then show ?case using * A-in by (metis bot-nat-0.extremum double-diff
sets.Diff sets.top sets-F-mono sets-le-imp-space-le space-F)
    }
next
  case (Union a)
    have a-in:  $a \subseteq \{0..\} \times \text{space } M$  for i using Union(1) sets.sets-into-space
sets-predictable-sigma space-predictable-sigma by metis
    hence A-in:  $A \subseteq \text{space } M$  for i using Union(4) by blast
    have snd  $x \in \text{snd } (a \cap (\{fst\ x\} \times \text{space } M))$  if  $x \in a$  for i x using that
a-in by fastforce
    hence a-i:  $a \cap i = (\bigcup j. \{j\} \times (\text{snd } (a \cap (\{j\} \times \text{space } M))))$  for i by force
    have A-i:  $A \cap i = \text{snd } (a \cap (\bigcup (\text{range } a) \cap (\{i\} \times \text{space } M)))$  for i unfolding
Union(4) using A-in by force
    have *:  $\text{snd } (a \cap (\{Suc\ i\} \times \text{space } M)) \in F\ i$   $\text{snd } (a \cap (\{0\} \times \text{space } M)) \in F\ 0$  for j using Union(2,3)[OF a-i] by auto
    {
      case 1
        have  $(\bigcup j. \text{snd } (a \cap (\{Suc\ i\} \times \text{space } M))) \in F\ i$  using * by fast
        moreover have  $(\bigcup j. \text{snd } (a \cap (\{Suc\ i\} \times \text{space } M))) = \text{snd } (a \cap (\bigcup (\text{range } a) \cap (\{Suc\ i\} \times \text{space } M)))$  by fast
        ultimately show ?case using A-i by metis
      next
        case 2
          have  $(\bigcup j. \text{snd } (a \cap (\{0\} \times \text{space } M))) \in F\ 0$  using * by fast
          moreover have  $(\bigcup j. \text{snd } (a \cap (\{0\} \times \text{space } M))) = \text{snd } (a \cap (\bigcup (\text{range } a) \cap (\{0\} \times \text{space } M)))$  by fast
          ultimately show ?case using A-i by metis
        }
    }
qed

```

This leads to the following useful fact.

lemma (in *nat-predictable-process*) *adapted-Suc*: *nat-adapted-process* *M F* ($\lambda i. X\ (Suc\ i)$)

proof (*unfold-locales*, *intro borel-measurableI*)

fix *S* :: 'b *set* **and** *i* **assume** *open-S*: *open* *S*

have $\{Suc\ i\} = \{i <.. Suc\ i\}$ **by** *fastforce*

hence $\{Suc\ i\} \times \text{space } M \in \Sigma_P$ **using** *space-F[symmetric, of i]* **unfolding** *sets-predictable-sigma* **by** (*intro sigma-sets.Basic*) *blast*

moreover **have** *case-prod* $X -' S \cap (UNIV \times \text{space } M) \in \Sigma_P$ **unfolding** *atLeast-0[symmetric]* **using** *open-S* **by** (*intro predictableD, simp add: borel-open*)

ultimately **have** *case-prod* $X -' S \cap (\{Suc\ i\} \times \text{space } M) \in \Sigma_P$ **unfolding** *sets-predictable-sigma* **using** *space-F* *sets.sets-into-space*

by (subst Times-Int-distrib1[of {Suc i} UNIV space M, simplified], subst
 inf.commute, subst Int-assoc[symmetric], subst Int-range-binary)
 (intro sigma-sets-Inter binary-in-sigma-sets, fast)+
 moreover have case-prod $X - 'S \cap (\{Suc\ i\} \times space\ M) = \{Suc\ i\} \times (X\ (Suc\ i) - 'S \cap space\ M)$ by (auto simp add: le-Suc-eq)
 moreover have $\dots = (\bigcup j. \{j\} \times (if\ j = Suc\ i\ then\ (X\ (Suc\ i) - 'S \cap space\ M)\ else\ \{\}))$ by (force split: if-splits)
 ultimately have $(\bigcup j. \{j\} \times (if\ j = Suc\ i\ then\ (X\ (Suc\ i) - 'S \cap space\ M)\ else\ \{\})) \in \Sigma_P$ by argo
 thus $X\ (Suc\ i) - 'S \cap space\ (F\ i) \in sets\ (F\ i)$ using sets-in-filtration[of $\lambda j. if\ j = Suc\ i\ then\ (X\ (Suc\ i) - 'S \cap space\ M)\ else\ \{\}$] space-F[OF zero-le] by
 presburger
 qed

theorem nat-predictable-process-iff: $nat\text{-predictable-process}\ M\ F\ X \longleftrightarrow nat\text{-adapted-process}\ M\ F\ (\lambda i. X\ (Suc\ i)) \wedge X\ 0 \in borel\text{-measurable}\ (F\ 0)$

proof (intro iffI)

assume asm: $nat\text{-adapted-process}\ M\ F\ (\lambda i. X\ (Suc\ i)) \wedge X\ 0 \in borel\text{-measurable}\ (F\ 0)$
 (F 0)

interpret nat-adapted-process $M\ F\ \lambda i. X\ (Suc\ i)$ **using** asm **by** blast

have $(\lambda(x, y). X\ x\ y) \in borel\text{-measurable}\ \Sigma_P$

proof (intro borel-measurableI)

fix $S :: 'b\ set$ **assume** open-S: open S

have $\{i\} \times (X\ i - 'S \cap space\ M) \in sets\ \Sigma_P$ **for** i

proof (cases i)

case 0

then show ?thesis **unfolding** sets-predictable-sigma

using measurable-sets[OF borel-open[OF open-S], of $X\ 0\ F\ 0$] asm **by** auto

next

case (Suc i)

have $\{Suc\ i\} = \{i <.. Suc\ i\}$ **by** fastforce

then show ?thesis **unfolding** sets-predictable-sigma

using measurable-sets[OF adapted borel-open[OF open-S], of i]

by (intro sigma-sets.Basic, auto simp add: Suc)

qed

moreover have $(\lambda(x, y). X\ x\ y) - 'S \cap space\ \Sigma_P = (\bigcup i. \{i\} \times (X\ i - 'S \cap space\ M))$ **by** fastforce

ultimately show $(\lambda(x, y). X\ x\ y) - 'S \cap space\ \Sigma_P \in sets\ \Sigma_P$ **by** simp

qed

thus $nat\text{-predictable-process}\ M\ F\ X$ **by** (unfold-locales)

next

assume asm: $nat\text{-predictable-process}\ M\ F\ X$

interpret $nat\text{-predictable-process}\ M\ F\ X$ **by** (rule asm)

show $nat\text{-adapted-process}\ M\ F\ (\lambda i. X\ (Suc\ i)) \wedge X\ 0 \in borel\text{-measurable}\ (F\ 0)$

using adapted-Suc **by** simp

qed

end

```

theory Martingale
  imports Stochastic-Process Conditional-Expectation-Banach
begin

```

11 Martingales

The following locales are necessary for defining martingales.

11.1 Additional Locale Definitions

```

locale sigma-finite-adapted-process = sigma-finite-filtered-measure  $M\ F\ t_0$  + adapted-process
 $M\ F\ t_0\ X$  for  $M\ F\ t_0\ X$ 

```

```

locale nat-sigma-finite-adapted-process = sigma-finite-adapted-process  $M\ F\ 0 :: nat$ 
 $X$  for  $M\ F\ X$ 

```

```

locale real-sigma-finite-adapted-process = sigma-finite-adapted-process  $M\ F\ 0 ::$ 
 $real\ X$  for  $M\ F\ X$ 

```

```

sublocale nat-sigma-finite-adapted-process  $\subseteq$  nat-sigma-finite-filtered-measure ..

```

```

sublocale real-sigma-finite-adapted-process  $\subseteq$  real-sigma-finite-filtered-measure ..

```

```

locale finite-adapted-process = finite-filtered-measure  $M\ F\ t_0$  + adapted-process  $M$ 
 $F\ t_0\ X$  for  $M\ F\ t_0\ X$ 

```

```

sublocale finite-adapted-process  $\subseteq$  sigma-finite-adapted-process ..

```

```

locale nat-finite-adapted-process = finite-adapted-process  $M\ F\ 0 :: nat\ X$  for  $M\ F\ X$ 

```

```

locale real-finite-adapted-process = finite-adapted-process  $M\ F\ 0 :: real\ X$  for  $M\ F\ X$ 

```

```

sublocale nat-finite-adapted-process  $\subseteq$  nat-sigma-finite-adapted-process ..

```

```

sublocale real-finite-adapted-process  $\subseteq$  real-sigma-finite-adapted-process ..

```

```

locale sigma-finite-adapted-process-order = sigma-finite-adapted-process  $M\ F\ t_0\ X$ 
for  $M\ F\ t_0$  and  $X :: - \Rightarrow - \Rightarrow - :: \{order-topology, ordered-real-vector\}$ 

```

```

locale nat-sigma-finite-adapted-process-order = sigma-finite-adapted-process-order
 $M\ F\ 0 :: nat\ X$  for  $M\ F\ X$ 

```

```

locale real-sigma-finite-adapted-process-order = sigma-finite-adapted-process-order
 $M\ F\ 0 :: real\ X$  for  $M\ F\ X$ 

```

```

sublocale nat-sigma-finite-adapted-process-order  $\subseteq$  nat-sigma-finite-adapted-process
..

```

```

sublocale real-sigma-finite-adapted-process-order  $\subseteq$  real-sigma-finite-adapted-process
..

locale finite-adapted-process-order = finite-adapted-process M F t0 X for M F t0
and X :: -  $\Rightarrow$  -  $\Rightarrow$  - :: {order-topology, ordered-real-vector}

locale nat-finite-adapted-process-order = finite-adapted-process-order M F 0 :: nat
X for M F X
locale real-finite-adapted-process-order = finite-adapted-process-order M F 0 :: real
X for M F X

sublocale nat-finite-adapted-process-order  $\subseteq$  nat-sigma-finite-adapted-process-order
..
sublocale real-finite-adapted-process-order  $\subseteq$  real-sigma-finite-adapted-process-order
..

locale sigma-finite-adapted-process-linorder = sigma-finite-adapted-process-order
M F t0 X for M F t0 and X :: -  $\Rightarrow$  -  $\Rightarrow$  - :: {linorder-topology}

locale nat-sigma-finite-adapted-process-linorder = sigma-finite-adapted-process-linorder
M F 0 :: nat X for M F X
locale real-sigma-finite-adapted-process-linorder = sigma-finite-adapted-process-linorder
M F 0 :: real X for M F X

sublocale nat-sigma-finite-adapted-process-linorder  $\subseteq$  nat-sigma-finite-adapted-process-order
..
sublocale real-sigma-finite-adapted-process-linorder  $\subseteq$  real-sigma-finite-adapted-process-order
..

locale finite-adapted-process-linorder = finite-adapted-process-order M F t0 X for
M F t0 and X :: -  $\Rightarrow$  -  $\Rightarrow$  - :: {linorder-topology}

locale nat-finite-adapted-process-linorder = finite-adapted-process-linorder M F 0
:: nat X for M F X
locale real-finite-adapted-process-linorder = finite-adapted-process-linorder M F 0
:: real X for M F X

sublocale nat-finite-adapted-process-linorder  $\subseteq$  nat-sigma-finite-adapted-process-linorder
..
sublocale real-finite-adapted-process-linorder  $\subseteq$  real-sigma-finite-adapted-process-linorder
..

```

11.2 Martingale

```

locale martingale = sigma-finite-adapted-process +
  assumes integrable:  $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$ 
  and martingale-property:  $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE \xi \text{ in } M. X i \xi =$ 

```

cond-exp $M (F i) (X j) \xi$

locale *martingale-order* = *martingale* $M F t_0 X$ **for** $M F t_0$ **and** $X :: - \Rightarrow - \Rightarrow -$
 $:: \{order-topology, ordered-real-vector\}$
locale *martingale-linorder* = *martingale* $M F t_0 X$ **for** $M F t_0$ **and** $X :: - \Rightarrow - \Rightarrow -$
 $- :: \{linorder-topology, ordered-real-vector\}$
sublocale *martingale-linorder* \subseteq *martingale-order* ..

lemma (in *sigma-finite-filtered-measure*) *martingale-const-fun*[intro]:
assumes *integrable* $M f f \in \text{borel-measurable } (F t_0)$
shows *martingale* $M F t_0 (\lambda-. f)$
using *assms sigma-finite-subalgebra.cond-exp-F-meas*[OF - *assms*(1), *THEN AE-symmetric*]
borel-measurable-mono
by (*unfold-locale*) *blast*+

lemma (in *sigma-finite-filtered-measure*) *martingale-cond-exp*[intro]:
assumes *integrable* $M f$
shows *martingale* $M F t_0 (\lambda i. \text{cond-exp } M (F i) f)$
using *sigma-finite-subalgebra.borel-measurable-cond-exp'* *borel-measurable-cond-exp*

by (*unfold-locale*) (*auto intro: sigma-finite-subalgebra.cond-exp-nested-subalg*[OF - *assms*]) *simp add: subalgebra-F subalgebra*)

corollary (in *sigma-finite-filtered-measure*) *martingale-zero*[intro]: *martingale* $M F t_0 (\lambda-. 0)$ **by** *fastforce*

corollary (in *finite-filtered-measure*) *martingale-const*[intro]: *martingale* $M F t_0 (\lambda-. c)$ **by** *fastforce*

11.3 Submartingale

locale *submartingale* = *sigma-finite-adapted-process-order* +
assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
and *submartingale-property*: $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE \xi \text{ in } M. X i \xi \leq$
cond-exp $M (F i) (X j) \xi$

locale *submartingale-linorder* = *submartingale* $M F t_0 X$ **for** $M F t_0$ **and** $X :: - \Rightarrow - \Rightarrow - :: \{linorder-topology\}$

sublocale *martingale-order* \subseteq *submartingale* **using** *martingale-property* **by** (*unfold-locale*)
(*force simp add: integrable*) +
sublocale *martingale-linorder* \subseteq *submartingale-linorder* ..

11.4 Supermartingale

locale *supermartingale* = *sigma-finite-adapted-process-order* +
assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
and *supermartingale-property*: $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE \xi \text{ in } M. X i \xi \geq$
cond-exp $M (F i) (X j) \xi$

locale *supermartingale-linorder* = *supermartingale* *M F t₀ X* **for** *M F t₀* **and** *X*
 $:: - \Rightarrow - \Rightarrow - :: \{ \text{linorder-topology} \}$

sublocale *martingale-order* \subseteq *supermartingale* **using** *martingale-property* **by** (*unfold-locales*)
(*force simp add: integrable*)
sublocale *martingale-linorder* \subseteq *supermartingale-linorder* ..

lemma *martingale-iff*:

shows *martingale M F t₀ X* \longleftrightarrow *submartingale M F t₀ X* \wedge *supermartingale M F t₀ X*

proof (*rule iffI*)

assume *asm: martingale M F t₀ X*

interpret *martingale-order M F t₀ X* **by** (*intro martingale-order.intro asm*)

show *submartingale M F t₀ X* \wedge *supermartingale M F t₀ X* **using** *submartingale-axioms supermartingale-axioms* **by** *blast*

next

assume *asm: submartingale M F t₀ X* \wedge *supermartingale M F t₀ X*

interpret *submartingale M F t₀ X* **by** (*simp add: asm*)

interpret *supermartingale M F t₀ X* **by** (*simp add: asm*)

show *martingale M F t₀ X* **using** *submartingale-property supermartingale-property*
by (*unfold-locales*) (*intro integrable, blast, force*)

qed

11.5 Martingale Lemmas

context *martingale*

begin

lemma *cond-exp-diff-eq-zero*:

assumes $t_0 \leq i \leq j$

shows $AE \xi \text{ in } M. \text{cond-exp } M (F i) (\lambda \xi. X j \xi - X i \xi) \xi = 0$

using *martingale-property[OF assms]* *assms*

sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i]

sigma-finite-subalgebra.cond-exp-diff[OF - integrable(1,1), of F i j i] **by**

fastforce

lemma *set-integral-eq*:

assumes $A \in F i \ t_0 \leq i \leq j$

shows *set-lebesgue-integral M A (X i)* = *set-lebesgue-integral M A (X j)*

proof –

interpret *sigma-finite-subalgebra M F i* **using** *assms(2)* **by** *blast*

have $\int x \in A. X i x \partial M = \int x \in A. \text{cond-exp } M (F i) (X j) x \partial M$ **using** *martingale-property[OF assms(2,3)] borel-measurable-cond-exp' assms subalgebra subalgebra-def* **by** (*intro set-lebesgue-integral-cong-AE[OF - random-variable]*) *fastforce* +

also have ... = $\int x \in A. X j x \partial M$ **using** *assms* **by** (*auto simp: integrable intro: cond-exp-set-integral[symmetric]*)

finally show *?thesis* .

qed

```

lemma scaleR-const[intro]:
  shows martingale  $M F t_0 (\lambda i x. c *_R X i x)$ 
proof –
  {
    fix  $i j :: 'b$  assume  $asm: t_0 \leq i i \leq j$ 
    interpret sigma-finite-subalgebra  $M F i$  using  $asm$  by blast
    have  $AE\ x\ in\ M. c *_R X i x = cond-exp\ M\ (F\ i)\ (\lambda x. c *_R X j x)\ x$  using
     $asm\ cond-exp-scaleR-right[OF\ integrable,\ of\ j,\ THEN\ AE-symmetric]\ martingale-property[OF\ asm]$  by force
  }
  thus ?thesis by (unfold-locales) (auto simp add: integrable martingale.integrable)
qed

```

```

lemma uminus[intro]:
  shows martingale  $M F t_0 (-\ X)$ 
  using scaleR-const[of  $-1$ ] by (force intro: back-subst[of martingale  $M F t_0$ ])

```

```

lemma add[intro]:
  assumes martingale  $M F t_0 Y$ 
  shows martingale  $M F t_0 (\lambda i \xi. X i \xi + Y i \xi)$ 
proof –
  interpret  $Y: martingale\ M\ F\ t_0\ Y$  by (rule assms)
  {
    fix  $i j :: 'b$  assume  $asm: t_0 \leq i i \leq j$ 
    hence  $AE\ \xi\ in\ M. X i \xi + Y i \xi = cond-exp\ M\ (F\ i)\ (\lambda x. X j x + Y j x)\ \xi$ 
    using sigma-finite-subalgebra.cond-exp-add[OF - integrable martingale.integrable][OF
    assms], of  $F\ i\ j\ j$ , THEN AE-symmetric]
    martingale-property[OF  $asm$ ] martingale.martingale-property[OF assms
     $asm$ ] by force
  }
  thus ?thesis using assms
  by (unfold-locales) (auto simp add: integrable martingale.integrable)
qed

```

```

lemma diff[intro]:
  assumes martingale  $M F t_0 Y$ 
  shows martingale  $M F t_0 (\lambda i x. X i x - Y i x)$ 
proof –
  interpret  $Y: martingale\ M\ F\ t_0\ Y$  by (rule assms)
  {
    fix  $i j :: 'b$  assume  $asm: t_0 \leq i i \leq j$ 
    hence  $AE\ \xi\ in\ M. X i \xi - Y i \xi = cond-exp\ M\ (F\ i)\ (\lambda x. X j x - Y j x)\ \xi$ 
    using sigma-finite-subalgebra.cond-exp-diff[OF - integrable martingale.integrable][OF
    assms], of  $F\ i\ j\ j$ , THEN AE-symmetric]
    martingale-property[OF  $asm$ ] martingale.martingale-property[OF assms
     $asm$ ] by fastforce
  }
  thus ?thesis using assms by (unfold-locales) (auto simp add: integrable martingale.integrable)

```

qed

end

lemma (in *sigma-finite-adapted-process*) *martingale-of-cond-exp-diff-eq-zero*:
assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
and *diff-zero*: $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE x \text{ in } M. \text{cond-exp } M (F i) (\lambda \xi. X j \xi - X i \xi) x = 0$
shows *martingale* $M F t_0 X$
proof
{
 fix $i j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
 thus $AE \xi \text{ in } M. X i \xi = \text{cond-exp } M (F i) (X j) \xi$
 using *diff-zero*[*OF asm*] *sigma-finite-subalgebra.cond-exp-diff*[*OF - integrable*(1,1),
of F i j i]
 sigma-finite-subalgebra.cond-exp-F-meas[*OF - integrable adapted, of i*] **by**
fastforce
}
qed (*intro integrable*)

lemma (in *sigma-finite-adapted-process*) *martingale-of-set-integral-eq*:
assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
and $\bigwedge A i j. t_0 \leq i \implies i \leq j \implies A \in F i \implies \text{set-lebesgue-integral } M A (X i) = \text{set-lebesgue-integral } M A (X j)$
shows *martingale* $M F t_0 X$
proof (*unfold-locale*)
fix $i j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
interpret *sigma-finite-subalgebra* $M F i$ **using** *asm* **by** *blast*
interpret *r*: *sigma-finite-measure restr-to-subalg* $M (F i)$ **by** (*simp add: sigma-fin-subalg*)
{
 fix A **assume** $A \in \text{restr-to-subalg } M (F i)$
 hence $*$: $A \in F i$ **using** *sets-restr-to-subalg subalgebra asm* **by** *blast*
 have $\text{set-lebesgue-integral } (\text{restr-to-subalg } M (F i)) A (X i) = \text{set-lebesgue-integral } M A (X i)$ **using** $*$ *subalg asm* **by** (*auto simp: set-lebesgue-integral-def intro: integral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator*)
 also have $\dots = \text{set-lebesgue-integral } M A (\text{cond-exp } M (F i) (X j))$ **using** $*$
 assms(2)[*OF asm*] *cond-exp-set-integral*[*OF integrable*] *asm* **by** *auto*
 finally have $\text{set-lebesgue-integral } (\text{restr-to-subalg } M (F i)) A (X i) = \text{set-lebesgue-integral } (\text{restr-to-subalg } M (F i)) A (\text{cond-exp } M (F i) (X j))$ **using** $*$ *subalg* **by** (*auto simp: set-lebesgue-integral-def intro!: integral-subalgebra2[symmetric] borel-measurable-scaleR borel-measurable-cond-exp borel-measurable-indicator*)
}
hence $AE \xi \text{ in } \text{restr-to-subalg } M (F i). X i \xi = \text{cond-exp } M (F i) (X j) \xi$ **using** *asm* **by** (*intro r.density-unique-banach, auto intro: integrable-in-subalg subalg borel-measurable-cond-exp integrable*)
thus $AE \xi \text{ in } M. X i \xi = \text{cond-exp } M (F i) (X j) \xi$ **using** *AE-restr-to-subalg*[*OF subalg*] **by** *blast*
qed (*simp add: integrable*)

11.6 Submartingale Lemmas

context *submartingale*

begin

lemma *cond-exp-diff-nonneg*:

assumes $t_0 \leq i \leq j$

shows $AE\ x\ in\ M. \ cond\text{-}exp\ M\ (F\ i)\ (\lambda\xi. X\ j\ \xi - X\ i\ \xi)\ x \geq 0$

using *submartingale-property*[*OF assms*] *assms sigma-finite-subalgebra.cond-exp-diff*[*OF - integrable(1,1), of - j i*] *sigma-finite-subalgebra.cond-exp-F-meas*[*OF - integrable adapted, of i*] **by** *fastforce*

lemma *add*[*intro*]:

assumes *submartingale* $M\ F\ t_0\ Y$

shows *submartingale* $M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi + Y\ i\ \xi)$

proof –

interpret $Y: \text{submartingale } M\ F\ t_0\ Y$ **by** (*rule assms*)

{

fix $i\ j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$

hence $AE\ \xi\ in\ M. X\ i\ \xi + Y\ i\ \xi \leq cond\text{-}exp\ M\ (F\ i)\ (\lambda x. X\ j\ x + Y\ j\ x)\ \xi$

using *sigma-finite-subalgebra.cond-exp-add*[*OF - integrable submartingale.integrable*[*OF assms*], *of F i j j*]

submartingale-property[*OF asm*] *submartingale.submartingale-property*[*OF assms asm*] *add-mono*[*of X i - - Y i -*] **by** *force*

}

thus *?thesis using assms by (unfold-locale) (auto simp add: borel-measurable-add random-variable adapted integrable Y.random-variable Y.adapted submartingale.integrable)*

qed

lemma *diff*[*intro*]:

assumes *supermartingale* $M\ F\ t_0\ Y$

shows *submartingale* $M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi - Y\ i\ \xi)$

proof –

interpret $Y: \text{supermartingale } M\ F\ t_0\ Y$ **by** (*rule assms*)

{

fix $i\ j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$

hence $AE\ \xi\ in\ M. X\ i\ \xi - Y\ i\ \xi \leq cond\text{-}exp\ M\ (F\ i)\ (\lambda x. X\ j\ x - Y\ j\ x)\ \xi$

using *sigma-finite-subalgebra.cond-exp-diff*[*OF - integrable supermartingale.integrable*[*OF assms*], *of F i j j*]

submartingale-property[*OF asm*] *supermartingale.supermartingale-property*[*OF assms asm*] *diff-mono*[*of X i - - Y i -*] **by** *force*

}

thus *?thesis using assms by (unfold-locale) (auto simp add: borel-measurable-diff random-variable adapted integrable Y.random-variable Y.adapted supermartingale.integrable)*

qed

lemma *scaleR-nonneg*:

assumes $c \geq 0$

shows *submartingale* $M F t_0 (\lambda i \xi. c *_R X i \xi)$
proof
 {
 fix $i j :: 'b$ **assume** $asm: t_0 \leq i i \leq j$
 thus $AE \xi$ *in* $M. c *_R X i \xi \leq cond-exp M (F i) (\lambda \xi. c *_R X j \xi) \xi$
 using *sigma-finite-subalgebra.cond-exp-scaleR-right*[*OF - integrable, of F i j c*] *submartingale-property*[*OF asm*] **by** (*fastforce intro!*: *scaleR-left-mono*[*OF - assms*])
 }
qed (*auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable random-variable adapted borel-measurable-const-scaleR*)

lemma *scaleR-le-zero*:
 assumes $c \leq 0$
 shows *supermartingale* $M F t_0 (\lambda i \xi. c *_R X i \xi)$
proof
 {
 fix $i j :: 'b$ **assume** $asm: t_0 \leq i i \leq j$
 thus $AE \xi$ *in* $M. c *_R X i \xi \geq cond-exp M (F i) (\lambda \xi. c *_R X j \xi) \xi$
 using *sigma-finite-subalgebra.cond-exp-scaleR-right*[*OF - integrable, of F i j c*] *submartingale-property*[*OF asm*]
 by (*fastforce intro!*: *scaleR-left-mono-neg*[*OF - assms*])
 }
qed (*auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable random-variable adapted borel-measurable-const-scaleR*)

lemma *uminus*[*intro*]:
 shows *supermartingale* $M F t_0 (- X)$
 unfolding *fun-Compl-def* **using** *scaleR-le-zero*[*of -1*] **by** *simp*

end

context *submartingale-linorder*
begin

lemma *set-integral-le*:
 assumes $A \in F i t_0 \leq i i \leq j$
 shows *set-lebesgue-integral* $M A (X i) \leq set-lebesgue-integral M A (X j)$
 using *submartingale-property*[*OF assms(2), of j*] *assms subalgebra*
 by (*subst sigma-finite-subalgebra.cond-exp-set-integral*[*OF - integrable assms(1), of j*])
 (*auto intro! scaleR-left-mono integral-mono-AE-banach integrable-mult-indicator integrable simp add: subalgebra-def set-lebesgue-integral-def*)

lemma *max*:
 assumes *submartingale-linorder* $M F t_0 Y$
 shows *submartingale-linorder* $M F t_0 (\lambda i \xi. max (X i \xi) (Y i \xi))$
proof (*unfold-locales*)
 interpret $Y: submartingale-linorder M F t_0 Y$ **by** (*rule assms*)

```

{
  fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
  have AE  $\xi$  in M.  $\max (X i \xi) (Y i \xi) \leq \max (\text{cond-exp } M (F i) (X j) \xi)$ 
  ( $\text{cond-exp } M (F i) (Y j) \xi$ ) using submartingale-property Y.submartingale-property
  asm unfolding max-def by fastforce
  thus AE  $\xi$  in M.  $\max (X i \xi) (Y i \xi) \leq \text{cond-exp } M (F i) (\lambda \xi. \max (X j \xi) (Y j \xi)) \xi$ 
  using sigma-finite-subalgebra.cond-exp-max[OF - integrable Y.integrable, of F i j]
  asm by (fast intro: order.trans)
}
show  $\bigwedge i. t_0 \leq i \implies (\lambda \xi. \max (X i \xi) (Y i \xi)) \in \text{borel-measurable } (F i) \bigwedge i. t_0 \leq i \implies \text{integrable } M (\lambda \xi. \max (X i \xi) (Y i \xi))$ 
by (force intro: Y.integrable integrable assms)+
qed

```

lemma max-0:

```

shows submartingale-linorder M F t_0 ( $\lambda i \xi. \max 0 (X i \xi)$ )
proof -
  interpret zero: martingale-linorder M F t_0  $\lambda - . 0$  by (force intro: martingale-linorder.intro martingale-order.intro)
  show ?thesis by (intro zero.max submartingale-linorder.intro submartingale-axioms)
qed

```

end

lemma (in sigma-finite-adapted-process-order) submartingale-of-cond-exp-diff-nonneg:

```

assumes integrable:  $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$ 
and diff-nonneg:  $\bigwedge i j. t_0 \leq i \implies i \leq j \implies \text{AE } x \text{ in } M. \text{cond-exp } M (F i) (\lambda \xi. X j \xi - X i \xi) x \geq 0$ 
shows submartingale M F t_0 X
proof (unfold-locale)
{
  fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
  thus AE  $\xi$  in M.  $X i \xi \leq \text{cond-exp } M (F i) (X j) \xi$ 
  using diff-nonneg[OF asm] sigma-finite-subalgebra.cond-exp-diff[OF - integrable(1,1), of F i j i]
  sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] by
  fastforce
}
qed (intro integrable)

```

lemma (in sigma-finite-adapted-process-linorder) submartingale-of-set-integral-le:

```

assumes integrable:  $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$ 
and  $\bigwedge A i j. t_0 \leq i \implies i \leq j \implies A \in F i \implies \text{set-lebesgue-integral } M A (X i) \leq \text{set-lebesgue-integral } M A (X j)$ 
shows submartingale M F t_0 X
proof (unfold-locale)
{
  fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
  interpret r: sigma-finite-measure restr-to-subalg M (F i) using asm sigma-finite-subalgebra.sigma-fin-subalg

```

by *blast*
 {
 fix A **assume** $A \in \text{restr-to-subalg } M (F i)$
 hence $*$: $A \in F i$ **using** *asm sets-restr-to-subalg subalgebra* **by** *blast*
 have $\text{set-lebesgue-integral } (\text{restr-to-subalg } M (F i)) A (X i) = \text{set-lebesgue-integral } M A (X i)$ **using** $*$ *asm subalgebra* **by** (*auto simp: set-lebesgue-integral-def intro: integral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator*)
 also have $\dots \leq \text{set-lebesgue-integral } M A (\text{cond-exp } M (F i) (X j))$ **using** $*$ *assms(2)[OF asm] asm sigma-finite-subalgebra.cond-exp-set-integral[OF - integrable]* **by** *fastforce*
 also have $\dots = \text{set-lebesgue-integral } (\text{restr-to-subalg } M (F i)) A (\text{cond-exp } M (F i) (X j))$ **using** $*$ *asm subalgebra* **by** (*auto simp: set-lebesgue-integral-def intro!: integral-subalgebra2[symmetric] borel-measurable-scaleR borel-measurable-cond-exp borel-measurable-indicator*)
 finally have $0 \leq \text{set-lebesgue-integral } (\text{restr-to-subalg } M (F i)) A (\lambda \xi. \text{cond-exp } M (F i) (X j) \xi - X i \xi)$ **using** $*$ *asm subalgebra* **by** (*subst set-integral-diff, auto simp add: set-integrable-def sets-restr-to-subalg intro!: integrable adapted integrable-in-subalg borel-measurable-scaleR borel-measurable-indicator borel-measurable-cond-exp integrable-mult-indicator*)
 }
 hence $AE \xi \text{ in } \text{restr-to-subalg } M (F i). 0 \leq \text{cond-exp } M (F i) (X j) \xi - X i \xi$
 by (*intro r.density-nonneg integrable-in-subalg asm subalgebra borel-measurable-diff borel-measurable-cond-exp adapted Bochner-Integration.integrable-diff integrable-cond-exp integrable*)
 thus $AE \xi \text{ in } M. X i \xi \leq \text{cond-exp } M (F i) (X j) \xi$ **using** *AE-restr-to-subalg[OF subalgebra] asm* **by** *simp*
 }
qed (*intro integrable*)

11.7 Supermartingale Lemmas

The following lemmas are exact duals of the ones for submartingales.

context *supermartingale*
begin

lemma *cond-exp-diff-nonneg*:
 assumes $t_0 \leq i \leq j$
 shows $AE x \text{ in } M. \text{cond-exp } M (F i) (\lambda \xi. X i \xi - X j \xi) x \geq 0$
 using *assms supermartingale-property[OF assms] sigma-finite-subalgebra.cond-exp-diff[OF - integrable(1,1), of F i i j]*
 sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] **by** *fastforce*

lemma *add[intro]*:
 assumes *supermartingale* $M F t_0 Y$
 shows *supermartingale* $M F t_0 (\lambda i \xi. X i \xi + Y i \xi)$
proof –
 interpret Y : *supermartingale* $M F t_0 Y$ **by** (*rule assms*)
 {

```

    fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
    hence AE  $\xi$  in  $M$ .  $X i \xi + Y i \xi \geq \text{cond-exp } M (F i) (\lambda x. X j x + Y j x) \xi$ 
    using sigma-finite-subalgebra.cond-exp-add[OF - integrable supermartingale.integrable[OF
assms], of F i j]
    supermartingale-property[OF asm] supermartingale.supermartingale-property[OF
assms asm] add-mono[of - X i - - Y i -] by force
  }
  thus ?thesis using assms by (unfold-locale) (auto simp add: borel-measurable-add
random-variable adapted integrable Y.random-variable Y.adapted supermartingale.integrable)

```

qed

```

lemma diff[intro]:
  assumes submartingale M F t_0 Y
  shows supermartingale M F t_0 ( $\lambda i \xi. X i \xi - Y i \xi$ )
proof -
  interpret Y: submartingale M F t_0 Y by (rule assms)
  {
    fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
    hence AE  $\xi$  in  $M$ .  $X i \xi - Y i \xi \geq \text{cond-exp } M (F i) (\lambda x. X j x - Y j x) \xi$ 
    using sigma-finite-subalgebra.cond-exp-diff[OF - integrable submartingale.integrable[OF
assms], of F i j j, unfolded fun-diff-def]
    supermartingale-property[OF asm] submartingale.submartingale-property[OF
assms asm] diff-mono[of - X i - Y i -] by force
  }
  thus ?thesis using assms by (unfold-locale) (auto simp add: borel-measurable-diff
random-variable adapted integrable Y.random-variable Y.adapted submartingale.integrable)

```

qed

```

lemma scaleR-nonneg:
  assumes  $c \geq 0$ 
  shows supermartingale M F t_0 ( $\lambda i \xi. c *_R X i \xi$ )
proof
  {
    fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
    thus AE  $\xi$  in  $M$ .  $c *_R X i \xi \geq \text{cond-exp } M (F i) (\lambda \xi. c *_R X j \xi) \xi$ 
    using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of F i
j c] supermartingale-property[OF asm] by (fastforce intro!: scaleR-left-mono[OF -
assms])
  }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)

```

```

lemma scaleR-le-zero:
  assumes  $c \leq 0$ 
  shows submartingale M F t_0 ( $\lambda i \xi. c *_R X i \xi$ )
proof
  {

```

```

    fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
    thus AE  $\xi$  in  $M$ .  $c *_R X i \xi \leq \text{cond-exp } M (F i) (\lambda \xi. c *_R X j \xi) \xi$ 
    using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of  $F i j c$ ]
    supermartingale-property[OF asm] by (fastforce intro!: scaleR-left-mono-neg[OF -
    assms])
  }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)

```

```

lemma uminus[intro]:
  shows submartingale  $M F t_0 (- X)$ 
  unfolding fun-Compl-def using scaleR-le-zero[of  $-1$ ] by simp

```

end

```

context supermartingale-linorder
begin

```

```

lemma set-integral-ge:
  assumes  $A \in F i t_0 \leq i \leq j$ 
  shows set-lebesgue-integral  $M A (X i) \geq \text{set-lebesgue-integral } M A (X j)$ 
  using supermartingale-property[OF assms(2), of  $j$ ] assms subalgebra
  by (subst sigma-finite-subalgebra.cond-exp-set-integral[OF - integrable assms(1),
  of  $j$ ])
  (auto intro!: scaleR-left-mono integral-mono-AE-banach integrable-mult-indicator
  integrable simp add: subalgebra-def set-lebesgue-integral-def)

```

```

lemma min:
  assumes supermartingale-linorder  $M F t_0 Y$ 
  shows supermartingale-linorder  $M F t_0 (\lambda i \xi. \min (X i \xi) (Y i \xi))$ 
proof (unfold-locales)
  interpret  $Y$ : supermartingale-linorder  $M F t_0 Y$  by (rule assms)
  {
    fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
    have AE  $\xi$  in  $M$ .  $\min (X i \xi) (Y i \xi) \geq \min (\text{cond-exp } M (F i) (X j) \xi) (\text{cond-exp } M (F i) (Y j) \xi)$ 
    using supermartingale-property  $Y$ .supermartingale-property asm
    unfolding min-def by fastforce
    thus AE  $\xi$  in  $M$ .  $\min (X i \xi) (Y i \xi) \geq \text{cond-exp } M (F i) (\lambda \xi. \min (X j \xi) (Y j \xi)) \xi$ 
    using sigma-finite-subalgebra.cond-exp-min[OF - integrable  $Y$ .integrable, of  $F i j j$ ] asm
    by (fast intro: order.trans)
  }
  show  $\bigwedge i. t_0 \leq i \implies (\lambda \xi. \min (X i \xi) (Y i \xi)) \in \text{borel-measurable } (F i) \bigwedge i. t_0 \leq i \implies \text{integrable } M (\lambda \xi. \min (X i \xi) (Y i \xi))$ 
  by (force intro:  $Y$ .integrable integrable assms)+
qed

```

```

lemma min-0:
  shows supermartingale-linorder  $M F t_0 (\lambda i \xi. \min 0 (X i \xi))$ 
proof -

```

interpret zero: martingale-linorder $M F t_0 \lambda \cdot . 0$ **by** (force intro: martingale-linorder.intro)
show ?thesis **by** (intro zero.min supermartingale-linorder.intro supermartingale-axioms)
qed

end

lemma (in sigma-finite-adapted-process-order) supermartingale-of-cond-exp-diff-le-zero:
assumes integrable: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
and diff-le-zero: $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE x \text{ in } M. \text{ cond-exp } M (F i)$
 $(\lambda \xi. X j \xi - X i \xi) x \leq 0$
shows supermartingale $M F t_0 X$
proof
{
 fix $i j :: 'b$ **assume** asm: $t_0 \leq i \leq j$
 thus $AE \xi \text{ in } M. X i \xi \geq \text{cond-exp } M (F i) (X j) \xi$
 using diff-le-zero[OF asm] sigma-finite-subalgebra.cond-exp-diff[OF - integrable(1,1), of $F i j i$]
 sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] **by**
 fastforce
}
qed (intro integrable)

lemma (in sigma-finite-adapted-process-linorder) supermartingale-of-set-integral-ge:
assumes integrable: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
and $\bigwedge A i j. t_0 \leq i \implies i \leq j \implies A \in F i \implies \text{set-lebesgue-integral } M A (X j) \leq \text{set-lebesgue-integral } M A (X i)$
shows supermartingale $M F t_0 X$
proof –
interpret -: adapted-process $M F t_0 -X$ **by** (rule uminus-adapted)
interpret uminus- X : sigma-finite-adapted-process-linorder $M F t_0 -X$..
note * = set-integral-uminus[unfolded set-integrable-def, OF integrable-mult-indicator[OF - integrable]]
have supermartingale $M F t_0 (-(-X))$
using ord-eq-le-trans[OF * ord-le-eq-trans[OF le-imp-neg-le[OF assms(2)] *[symmetric]]]
subalgebra
by (intro submartingale.uminus uminus- X .submartingale-of-set-integral-le)
(clarsimp simp add: fun-Compl-def subalgebra-def integrable | fastforce)+
thus ?thesis **unfolding** fun-Compl-def **by** simp
qed

11.8 Discrete Time Martingales

locale nat-martingale = martingale $M F 0 :: \text{nat } X$ **for** $M F X$
locale nat-submartingale = submartingale $M F 0 :: \text{nat } X$ **for** $M F X$
locale nat-supermartingale = supermartingale $M F 0 :: \text{nat } X$ **for** $M F X$

locale nat-submartingale-linorder = submartingale-linorder $M F 0 :: \text{nat } X$ **for** M

```

F X
locale nat-supermartingale-linorder = supermartingale-linorder M F 0 :: nat X
for M F X

sublocale nat-submartingale-linorder  $\subseteq$  nat-submartingale ..
sublocale nat-supermartingale-linorder  $\subseteq$  nat-supermartingale ..

lemma (in nat-martingale) predictable-const:
  assumes nat-predictable-process M F X
  shows AE  $\xi$  in M. X i  $\xi$  = X j  $\xi$ 
proof -
  have *: AE  $\xi$  in M. X i  $\xi$  = X 0  $\xi$  for i
  proof (induction i)
    case 0
    then show ?case by (simp add: bot-nat-def)
  next
    case (Suc i)
    interpret S: nat-adapted-process M F  $\lambda i$ . X (Suc i) by (intro nat-predictable-process.adapted-Suc
  asms)
    show ?case using Suc S.adapted[of i] martingale-property[OF - le-SucI, of i]
  sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable, of F i Suc i] by fastforce
  qed
  show ?thesis using *[of i] *[of j] by force
qed

lemma (in nat-sigma-finite-adapted-process) martingale-of-set-integral-eq-Suc:
  assumes integrable:  $\bigwedge i$ . integrable M (X i)
  and  $\bigwedge A$  i. A  $\in F$  i  $\implies$  set-lebesgue-integral M A (X i) = set-lebesgue-integral
  M A (X (Suc i))
  shows nat-martingale M F X
proof (intro nat-martingale.intro martingale-of-set-integral-eq)
  fix i j A assume asm: i  $\leq$  j A  $\in$  sets (F i)
  show set-lebesgue-integral M A (X i) = set-lebesgue-integral M A (X j) using
  asm
  proof (induction j - i arbitrary: i j)
    case 0
    then show ?case using asm by simp
  next
    case (Suc n)
    hence *: n = j - Suc i by linarith
    have Suc i  $\leq$  j using Suc(2,3) by linarith
    thus ?case using sets-F-mono[OF - le-SucI] Suc(4) Suc(1)[OF *] by (auto
  intro: asms(2)[THEN trans])
  qed
qed (simp add: integrable)

lemma (in nat-sigma-finite-adapted-process) martingale-nat:
  assumes integrable:  $\bigwedge i$ . integrable M (X i)
  and  $\bigwedge i$ . AE  $\xi$  in M. X i  $\xi$  = cond-exp M (F i) (X (Suc i))  $\xi$ 

```

shows *nat-martingale* $M F X$
proof (*unfold-locales*)
fix $i j :: \text{nat}$ **assume** $\text{asm}: i \leq j$
show $AE \xi \text{ in } M. X i \xi = \text{cond-exp } M (F i) (X j) \xi$ **using** asm
proof (*induction* $j - i$ *arbitrary: i j*)
case 0
hence $j = i$ **by** *simp*
thus ?*case* **using** *sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, THEN AE-symmetric]* **by** *blast*
next
case (*Suc n*)
have $j: j = \text{Suc } (n + i)$ **using** *Suc* **by** *linarith*
have $n: n = n + i - i$ **using** *Suc* **by** *linarith*
have *: $AE \xi \text{ in } M. \text{cond-exp } M (F (n + i)) (X j) \xi = X (n + i) \xi$ **unfolding**
 j **using** *assms(2)[THEN AE-symmetric]* **by** *blast*
have $AE \xi \text{ in } M. \text{cond-exp } M (F i) (X j) \xi = \text{cond-exp } M (F i) (\text{cond-exp } M$
 $(F (n + i)) (X j)) \xi$ **by** (*intro cond-exp-nested-subalg integrable subalg, simp add:*
subalgebra-def sets-F-mono)
hence $AE \xi \text{ in } M. \text{cond-exp } M (F i) (X j) \xi = \text{cond-exp } M (F i) (X (n + i))$
 ξ **using** *cond-exp-cong-AE[OF integrable-cond-exp integrable *]* **by** *force*
thus ?*case* **using** *Suc(1)[OF n]* **by** *fastforce*
qed
qed (*simp add: integrable*)

lemma (*in nat-sigma-finite-adapted-process*) *martingale-of-cond-exp-diff-Suc-eq-zero*:
assumes *integrable: $\bigwedge i. \text{integrable } M (X i)$*
and $\bigwedge i. AE \xi \text{ in } M. \text{cond-exp } M (F i) (\lambda \xi. X (\text{Suc } i) \xi - X i \xi) \xi = 0$
shows *nat-martingale* $M F X$
proof (*intro martingale-nat integrable*)
fix i
show $AE \xi \text{ in } M. X i \xi = \text{cond-exp } M (F i) (X (\text{Suc } i)) \xi$ **using** *cond-exp-diff[OF*
integrable(1,1), of i Suc i] *cond-exp-F-meas[OF integrable adapted, of i]* *assms(2)[of*
 $i]$ **by** *fastforce*
qed

11.9 Discrete Time Submartingales

lemma (*in nat-submartingale*) *predictable-mono*:
assumes *nat-predictable-process* $M F X i \leq j$
shows $AE \xi \text{ in } M. X i \xi \leq X j \xi$
using *assms(2)*
proof (*induction* $j - i$ *arbitrary: i j*)
case 0
then show ?*case* **by** *simp*
next
case (*Suc n*)
hence *: $n = j - \text{Suc } i$ **by** *linarith*
interpret $S: \text{nat-adapted-process } M F \lambda i. X (\text{Suc } i)$ **by** (*intro nat-predictable-process.adapted-Suc*
assms)

have $Suc\ i \leq j$ **using** $Suc(2,3)$ **by** *linarith*
thus $?case$ **using** $Suc(1)[OF\ *]$ $S.adapted[of\ i]$ *submartingale-property* $[OF -$
 $le-SucI, of\ i]$ *sigma-finite-subalgebra.cond-exp-F-meas* $[OF - integrable, of\ F\ i\ Suc$
 $i]$ **by** *fastforce*
qed

lemma (*in nat-sigma-finite-adapted-process-linorder*) *submartingale-of-set-integral-le-Suc*:
assumes *integrable*: $\bigwedge i. integrable\ M\ (X\ i)$
and $\bigwedge A\ i. A \in F\ i \implies set-lebesgue-integral\ M\ A\ (X\ i) \leq set-lebesgue-integral$
 $M\ A\ (X\ (Suc\ i))$
shows *nat-submartingale* $M\ F\ X$
proof (*intro nat-submartingale.intro submartingale-of-set-integral-le*)
fix $i\ j\ A$ **assume** *asm*: $i \leq j\ A \in sets\ (F\ i)$
show $set-lebesgue-integral\ M\ A\ (X\ i) \leq set-lebesgue-integral\ M\ A\ (X\ j)$ **using**
asm
proof (*induction j - i arbitrary: i j*)
case 0
then show $?case$ **using** *asm* **by** *simp*
next
case ($Suc\ n$)
hence $*$: $n = j - Suc\ i$ **by** *linarith*
have $Suc\ i \leq j$ **using** $Suc(2,3)$ **by** *linarith*
thus $?case$ **using** *sets-F-mono* $[OF - le-SucI]$ $Suc(4)$ $Suc(1)[OF\ *]$ **by** (*auto*
intro: assms(2)[THEN order-trans])
qed
qed (*simp add: integrable*)

lemma (*in nat-sigma-finite-adapted-process-linorder*) *submartingale-nat*:
assumes *integrable*: $\bigwedge i. integrable\ M\ (X\ i)$
and $\bigwedge i. AE\ \xi\ in\ M. X\ i\ \xi \leq cond-exp\ M\ (F\ i)\ (X\ (Suc\ i))\ \xi$
shows *nat-submartingale* $M\ F\ X$
using *subalg integrable assms(2)*
by (*intro submartingale-of-set-integral-le-Suc ord-le-eq-trans* $[OF\ set-integral-mono-AE-banach$
 $cond-exp-set-integral[symmetric]], simp$)
(meson in-mono integrable-mult-indicator set-integrable-def subalgebra-def, me-
son integrable-cond-exp in-mono integrable-mult-indicator set-integrable-def subal-
gebra-def, fast+)

lemma (*in nat-sigma-finite-adapted-process-linorder*) *submartingale-of-cond-exp-diff-Suc-nonneg*:
assumes *integrable*: $\bigwedge i. integrable\ M\ (X\ i)$
and $\bigwedge i. AE\ \xi\ in\ M. cond-exp\ M\ (F\ i)\ (\lambda\xi. X\ (Suc\ i)\ \xi - X\ i\ \xi)\ \xi \geq 0$
shows *nat-submartingale* $M\ F\ X$
proof (*intro submartingale-nat integrable*)
fix i
show $AE\ \xi\ in\ M. X\ i\ \xi \leq cond-exp\ M\ (F\ i)\ (X\ (Suc\ i))\ \xi$ **using** *cond-exp-diff* $[OF$
 $integrable(1,1), of\ i\ Suc\ i\ i]$ *cond-exp-F-meas* $[OF\ integrable\ adapted, of\ i]$ *assms(2)[of*
 $i]$ **by** *fastforce*
qed

lemma (in *nat-submartingale-linorder*) *partial-sum-scaleR*:
assumes *nat-adapted-process* $M F C \wedge i. AE \xi \text{ in } M. 0 \leq C i \xi \wedge i. AE \xi \text{ in } M. C i \xi \leq R$
shows *nat-submartingale* $M F (\lambda n \xi. \sum i < n. C i \xi *_{\mathcal{R}} (X (Suc i) \xi - X i \xi))$
proof –
interpret C : *nat-adapted-process* $M F C$ **by** (*rule assms*)
interpret C' : *nat-adapted-process* $M F \lambda i \xi. C (i - 1) \xi *_{\mathcal{R}} (X i \xi - X (i - 1) \xi)$ **by** (*intro nat-adapted-process.intro adapted-process.scaleR-right-adapted adapted-process.diff-adapted, unfold-locales*) (*auto intro: adaptedD C.adaptedD*) +
interpret C'' : *nat-adapted-process* $M F \lambda n \xi. \sum i < n. C i \xi *_{\mathcal{R}} (X (Suc i) \xi - X i \xi)$ **by** (*rule C'.partial-sum-Suc-adapted[unfolded diff-Suc-1]*)
interpret S : *nat-sigma-finite-adapted-process-linorder* $M F (\lambda n \xi. \sum i < n. C i \xi *_{\mathcal{R}} (X (Suc i) \xi - X i \xi))$..
have *integrable* $M (\lambda x. C i x *_{\mathcal{R}} (X (Suc i) x - X i x))$ **for** i **using** *assms*(2,3)[*of i*] **by** (*intro Bochner-Integration.integrable-bound[OF integrable-scaleR-right, OF Bochner-Integration.integrable-diff, OF integrable(1,1), of R Suc i i]*) (*auto simp add: mult-mono*)
moreover have $AE \xi \text{ in } M. 0 \leq \text{cond-exp } M (F i) (\lambda \xi. (\sum i < Suc i. C i \xi *_{\mathcal{R}} (X (Suc i) \xi - X i \xi)) - (\sum i < i. C i \xi *_{\mathcal{R}} (X (Suc i) \xi - X i \xi))) \xi$ **for** i
using *sigma-finite-subalgebra.cond-exp-measurable-scaleR*[*OF - calculation - C.adapted, of i*]
cond-exp-diff-nonneg[*OF - le-SucI, OF - order.refl, of i*] *assms*(2,3)[*of i*]
by (*fastforce simp add: scaleR-nonneg-nonneg integrable*)
ultimately show ?thesis **by** (*intro S.submartingale-of-cond-exp-diff-Suc-nonneg Bochner-Integration.integrable-sum, blast+*)
qed

lemma (in *nat-submartingale-linorder*) *partial-sum-scaleR'*:
assumes *nat-predictable-process* $M F C \wedge i. AE \xi \text{ in } M. 0 \leq C i \xi \wedge i. AE \xi \text{ in } M. C i \xi \leq R$
shows *nat-submartingale* $M F (\lambda n \xi. \sum i < n. C (Suc i) \xi *_{\mathcal{R}} (X (Suc i) \xi - X i \xi))$
proof –
interpret C : *nat-predictable-process* $M F C$ **by** (*rule assms*)
interpret $Suc\text{-}C$: *nat-adapted-process* $M F \lambda i. C (Suc i)$ **using** $C.adapted\text{-}Suc$.
show ?thesis **by** (*intro partial-sum-scaleR[of - R] assms*) (*intro-locales*)
qed

11.10 Discrete Time Supermartingales

lemma (in *nat-supermartingale*) *predictable-mono*:

assumes *nat-predictable-process* $M F X i \leq j$

shows $AE \xi \text{ in } M. X i \xi \geq X j \xi$

using *assms*(2)

proof (*induction j - i arbitrary: i j*)

case 0

then show ?case **by** *simp*

next

case ($Suc n$)

hence *: $n = j - \text{Suc } i$ by *linarith*
 interpret S : *nat-adapted-process* $M F \lambda i. X (\text{Suc } i)$ by (intro *nat-predictable-process.adapted-Suc* *assms*)
 have $\text{Suc } i \leq j$ using $\text{Suc}(2,3)$ by *linarith*
 thus ?case using $\text{Suc}(1)[OF *]$ $S.\text{adapted}[of i]$ *supermartingale-property*[$OF -$ *le-SucI*, *of i*] *sigma-finite-subalgebra.cond-exp-F-meas*[$OF -$ *integrable*, *of F i Suc i*] by *fastforce*
 qed

lemma (in *nat-sigma-finite-adapted-process-linorder*) *supermartingale-of-set-integral-ge-Suc*:
 assumes *integrable*: $\bigwedge i. \text{integrable } M (X i)$
 and $\bigwedge A i. A \in F i \implies \text{set-lebesgue-integral } M A (X i) \geq \text{set-lebesgue-integral } M A (X (\text{Suc } i))$
 shows *nat-supermartingale* $M F X$
proof –
 interpret -: *adapted-process* $M F 0 -X$ by (rule *uminus-adapted*)
 interpret *uminus-X*: *nat-sigma-finite-adapted-process-linorder* $M F -X$..
 note * = *set-integral-uminus*[*unfolded set-integrable-def*, *OF integrable-mult-indicator*[$OF -$ *integrable*]]
 have *nat-supermartingale* $M F (-(-X))$
 using *ord-eq-le-trans*[$OF * \text{ord-le-eq-trans}[OF \text{le-imp-neg-le}[OF \text{assms}(2)]] * [\text{symmetric}]]$ *subalgebra*
 by (intro *nat-supermartingale.intro submartingale.uminus nat-submartingale.axioms uminus-X.submartingale-of-set-integral-le-Suc*)
 (clarsimp simp add: *fun-Compl-def subalgebra-def integrable* | *fastforce*) +
 thus ?thesis **unfolding** *fun-Compl-def* by *simp*
 qed

lemma (in *nat-sigma-finite-adapted-process-linorder*) *supermartingale-nat*:
 assumes *integrable*: $\bigwedge i. \text{integrable } M (X i)$
 and $\bigwedge i. AE \xi \text{ in } M. X i \xi \geq \text{cond-exp } M (F i) (X (\text{Suc } i)) \xi$
 shows *nat-supermartingale* $M F X$
proof –
 interpret -: *adapted-process* $M F 0 -X$ by (rule *uminus-adapted*)
 interpret *uminus-X*: *nat-sigma-finite-adapted-process-linorder* $M F -X$..
 have $AE \xi \text{ in } M. -X i \xi \leq \text{cond-exp } M (F i) (\lambda x. -X (\text{Suc } i) x) \xi$ for i using *assms*(2) *cond-exp-uminus*[$OF \text{integrable}$, *of i Suc i*] by *force*
 hence *nat-supermartingale* $M F (-(-X))$ by (intro *nat-supermartingale.intro submartingale.uminus nat-submartingale.axioms uminus-X.submartingale-nat*) (auto simp add: *fun-Compl-def integrable*)
 thus ?thesis **unfolding** *fun-Compl-def* by *simp*
 qed

lemma (in *nat-sigma-finite-adapted-process-linorder*) *supermartingale-of-cond-exp-diff-Suc-le-zero*:
 assumes *integrable*: $\bigwedge i. \text{integrable } M (X i)$
 and $\bigwedge i. AE \xi \text{ in } M. \text{cond-exp } M (F i) (\lambda \xi. X (\text{Suc } i) \xi - X i \xi) \xi \leq 0$
 shows *nat-supermartingale* $M F X$
proof (intro *supermartingale-nat integrable*)
 fix i

```

show  $\text{AE } \xi \text{ in } M. X \ i \ \xi \geq \text{cond-exp } M \ (F \ i) \ (X \ (\text{Suc } i)) \ \xi$  using cond-exp-diff[OF
integrable(1,1), of i Suc i i] cond-exp-F-meas[OF integrable adapted, of i] assms(2)[of
i] by fastforce
qed

end

```