

On the Formalization of Martingales

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Abstract

In the scope of this project, we present a formalization of martingales in arbitrary Banach spaces using Isabelle/HOL.

The current formalization of conditional expectation in the Isabelle library is limited to real-valued functions. To overcome this limitation, we extend the construction of conditional expectation to general Banach spaces, employing an approach similar to the one described in [1]. We use measure theoretic arguments to construct the conditional expectation using suitable limits of simple functions.

Subsequently, we define stochastic processes and introduce the concepts of adapted, progressively measurable and predictable processes using suitable locale definitions¹. We show the relation

$$\text{adapted} \supseteq \text{progressive} \supseteq \text{predictable}$$

Furthermore, we show that progressive measurability and adaptedness are equivalent when the indexing set is discrete. We pay special attention to predictable processes in discrete-time, showing that $(X_n)_{n \in \mathbb{N}}$ is predictable if and only if $(X_{n+1})_{n \in \mathbb{N}}$ is adapted.

Moving forward, we rigorously define martingales, submartingales, and supermartingales, presenting their first consequences and corollaries². Discrete-time martingales are given special attention in the formalization. In every step of our formalization, we make extensive use of the powerful locale system of Isabelle.

The formalization further contributes by generalizing concepts in Bochner integration by extending their application from the real numbers to arbitrary Banach spaces equipped with a second-countable topology. Induction schemes for integrable simple functions on Banach spaces are introduced, accommodating various scenarios with or without a real vector ordering³. Specifically, we formalize a powerful result called the “Averaging Theorem”[3] which allows us to show that densities are unique in Banach spaces.

In depth information on the formalization and the proofs of the individual theorems can be found in [2].

¹Martingale.Stochastic_Process

²Martingale.Martingale

³Martingale.Bochner_Integration_Addendum

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```

theory Measure-Space-Supplement
  imports HOL-Analysis.Measure-Space
begin

```

1 Supplementary Lemmas for Measure Spaces

1.1 σ -Algebra Generated by a Family of Functions

definition *family-vimage-algebra* :: '*a* set \Rightarrow ('*a* \Rightarrow '*b*) set \Rightarrow '*b* measure \Rightarrow '*a* measure **where**
family-vimage-algebra Ω *S* *M* \equiv *sigma* Ω ($\bigcup f \in S. \{f - 'A \cap \Omega \mid A. A \in M\}$)

For singleton *S*, i.e. $S = \{f\}$ for some *f*, the definition simplifies to that of *vimage-algebra*.

lemma *family-vimage-algebra-singleton*: *family-vimage-algebra* Ω $\{f\}$ *M* = *vimage-algebra* Ω *f* *M* **unfolding** *family-vimage-algebra-def* *vimage-algebra-def* **by** *simp*

lemma

shows *sets-family-vimage-algebra*: *sets* (*family-vimage-algebra* Ω *S* *M*) = *sigma-sets* Ω ($\bigcup f \in S. \{f - 'A \cap \Omega \mid A. A \in M\}$)
and *space-family-vimage-algebra*[*simp*]: *space* (*family-vimage-algebra* Ω *S* *M*) = Ω
by (*auto simp add: family-vimage-algebra-def sets-measure-of-conv space-measure-of-conv*)

lemma *measurable-family-vimage-algebra*:

assumes $f \in S$ $f \in \Omega \rightarrow \text{space } M$
shows $f \in \text{family-vimage-algebra } \Omega$ *S* *M* $\rightarrow_M M$
using *assms* **by** (*intro measurableI, auto simp add: sets-family-vimage-algebra*)

lemma *measurable-family-vimage-algebra-singleton*:

assumes $f \in \Omega \rightarrow \text{space } M$
shows $f \in \text{family-vimage-algebra } \Omega$ $\{f\}$ *M* $\rightarrow_M M$
using *assms* *measurable-family-vimage-algebra* **by** *blast*

A collection of functions are measurable with respect to some σ -algebra *N*, if and only if the σ -algebra they generate is contained in *N*.

lemma *measurable-family-iff-sets*:

shows ($S \subseteq N \rightarrow_M M$) \longleftrightarrow $S \subseteq \text{space } N \rightarrow \text{space } M \wedge \text{family-vimage-algebra} (\text{space } N) S M \subseteq N$

proof (*standard, goal-cases*)

case 1

hence *subset*: $S \subseteq \text{space } N \rightarrow \text{space } M$ **using** *measurable-space* **by** *fast*

have $\{f - 'A \cap \text{space } N \mid A. A \in M\} \subseteq N$ **if** $f \in S$ **for** *f* **using** *measurable-iff-sets*[*unfolded family-vimage-algebra-singleton[symmetric], of f*] 1 *subset that*
by (*fastforce simp add: sets-family-vimage-algebra*)

then show ?*case* **unfolding** *sets-family-vimage-algebra* **using** *sets.sigma-algebra-axioms*
by (*simp add: subset, intro sigma-algebra.sigma-sets-subset, blast+*)

```

next
  case 2
  hence subset:  $S \subseteq \text{space } N \rightarrow \text{space } M$  by simp
  show ?case
  proof (standard, goal-cases)
    case (1 x)
    have family-vimage-algebra ( $\text{space } N$ )  $\{x\}$   $M \subseteq N$  by (metis (no-types, lifting)
1 2 sets-family-vimage-algebra SUP-le-iff sigma-sets-le-sets-iff singletonD)
    thus ?case using measurable-iff-sets[unfolded family-vimage-algebra-singleton[symmetric]]
subset[THEN subsetD, OF 1] by fast
  qed
qed

lemma family-vimage-algebra-diff:
  shows family-vimage-algebra  $\Omega$   $S$   $M = \text{sigma } \Omega$  (sets (family-vimage-algebra  $\Omega$ 
( $S - I$ )  $M$ )  $\cup$  family-vimage-algebra  $\Omega$  ( $S \cap I$ )  $M$ )
  using sets.space-closed space-measure-of-conv
  unfolding family-vimage-algebra-def sets-family-vimage-algebra
  by (intro sigma-eqI, blast, fastforce)
    (intro sigma-sets-eqI, blast, simp add: sets-measure-of-conv split: if-splits,
meson Diff-subset Sup-subset-mono in-mono inf-sup-ord(1) sigma-sets-subseteq
subset-image-iff, fastforce+)

end
theory Elementary-Metric-Spaces-Supplement
  imports HOL-Analysis.Elementary-Metric-Spaces
begin

```

2 Supplementary Lemmas for Elementary Metric Spaces

2.1 Diameter Lemma

```

lemma diameter-comp-strict-mono:
  fixes  $s :: \text{nat} \Rightarrow 'a :: \text{metric-space}$ 
  assumes strict-mono  $r$  bounded  $\{s\ i \mid i. r\ n \leq i\}$ 
  shows diameter  $\{s\ (r\ i) \mid i. n \leq i\} \leq \text{diameter } \{s\ i \mid i. r\ n \leq i\}$ 
proof (rule diameter-subset)
  show  $\{s\ (r\ i) \mid i. n \leq i\} \subseteq \{s\ i \mid i. r\ n \leq i\}$  using assms(1) monotoneD
strict-mono-mono by fastforce
qed (intro assms(2))

lemma diameter-bounded-bound':
  fixes  $S :: 'a :: \text{metric-space}$  set
  assumes  $S$ : bdd-above (case-prod dist ' ( $S \times S$ ))  $x \in S\ y \in S$ 
  shows dist  $x\ y \leq \text{diameter } S$ 
proof -
  have  $(x,y) \in S \times S$  using  $S$  by auto
  then have dist  $x\ y \leq (\text{SUP } (x,y) \in S \times S. \text{dist } x\ y)$  by (rule cSUP-upper2[OF

```

```

assms(1)]) simp
  with  $\langle x \in S \rangle$  show ?thesis by (auto simp: diameter-def)
qed

```

```

lemma bounded-imp-dist-bounded:
  assumes bounded (range s)
  shows bounded (( $\lambda(i, j). \text{dist } (s \ i) (s \ j)$ ) ' ( $\{n..\} \times \{n..\}$ ))
  using bounded-dist-comp[OF bounded-fst bounded-snd, OF bounded-Times(1,1)[OF
assms(1,1)]] by (rule bounded-subset, force)

```

A sequence is Cauchy, if and only if it is bounded and it's diameter tends to zero. The diameter is well-defined only if the sequence is bounded.

```

lemma cauchy-iff-diameter-tends-to-zero-and-bounded:
  fixes s :: nat  $\Rightarrow$  'a :: metric-space
  shows Cauchy s  $\longleftrightarrow$  (( $\lambda n. \text{diameter } \{s \ i \mid i. i \geq n\}$ )  $\longrightarrow$  0  $\wedge$  bounded (range s))
proof -
  have {s i | i. N  $\leq$  i}  $\neq$  {} for N by blast
  hence diameter-SUP: diameter {s i | i. N  $\leq$  i} = (SUP (i, j)  $\in$  {N..}  $\times$  {N..}.
dist (s i) (s j)) for N unfolding diameter-def by (auto intro!: arg-cong[of - - Sup])
  show ?thesis
  proof (intro iffI)
    assume asm: Cauchy s
    have  $\exists N. \forall n \geq N. \text{norm } (\text{diameter } \{s \ i \mid i. n \leq i\}) < e$  if e-pos: e > 0 for e
    proof -
      obtain N where dist-less: dist (s n) (s m) < (1/2) * e if n  $\geq$  N m  $\geq$  N
      for n m using asm e-pos by (metis Cauchy-def mult-pos-pos zero-less-divide-iff
zero-less-numeral zero-less-one)
      {
        fix r assume r  $\geq$  N
        hence dist (s n) (s m) < (1/2) * e if n  $\geq$  r m  $\geq$  r for n m using dist-less
        that by simp
        hence (SUP (i, j)  $\in$  {r..}  $\times$  {r..}. dist (s i) (s j))  $\leq$  (1/2) * e by (intro
cSup-least) fastforce+
        also have ... < e using e-pos by simp
        finally have diameter {s i | i. r  $\leq$  i} < e using diameter-SUP by presburger
      }
      moreover have diameter {s i | i. r  $\leq$  i}  $\geq$  0 for r unfolding diameter-SUP
      using bounded-imp-dist-bounded[OF cauchy-imp-bounded, THEN bounded-imp-bdd-above,
OF asm] by (intro cSup-upper2, auto)
      ultimately show ?thesis by auto
    qed
    thus ( $\lambda n. \text{diameter } \{s \ i \mid i. n \leq i\}$ )  $\longrightarrow$  0  $\wedge$  bounded (range s) using
cauchy-imp-bounded[OF asm] by (simp add: LIMSEQ-iff)
  next
    assume asm: ( $\lambda n. \text{diameter } \{s \ i \mid i. n \leq i\}$ )  $\longrightarrow$  0  $\wedge$  bounded (range s)
    have  $\exists N. \forall n \geq N. \forall m \geq N. \text{dist } (s \ n) (s \ m) < e$  if e-pos: e > 0 for e
    proof -
      obtain N where diam-less: diameter {s i | i. r  $\leq$  i} < e if r  $\geq$  N for r

```

```

using LIMSEQ-D asm(1) e-pos by fastforce
{
  fix n m assume  $n \geq N$   $m \geq N$ 
  hence  $\text{dist } (s\ n) (s\ m) < e$  using cSUP-lessD[OF bounded-imp-dist-bounded[THEN
bounded-imp-bdd-above], OF - diam-less[unfolded diameter-SUP]] asm by auto
}
thus ?thesis by blast
qed
then show Cauchy s by (simp add: Cauchy-def)
qed
qed
end

```

```

theory Bochner-Integration-Supplement
imports HOL-Analysis.Bochner-Integration HOL-Analysis.Set-Integral Elementary-Metric-Spaces-Supplement
begin

```

3 Supplementary Lemmas for Bochner Integration

3.1 Integrable Simple Functions

We restate some basic results concerning Bochner-integrable functions.

```

lemma integrable-implies-simple-function-sequence:
  fixes  $f :: 'a \Rightarrow 'b :: \{\text{banach, second-countable-topology}\}$ 
  assumes integrable M f
  obtains s where  $\bigwedge i. \text{simple-function } M (s\ i)$ 
    and  $\bigwedge i. \text{emeasure } M \{y \in \text{space } M. s\ i\ y \neq 0\} \neq \infty$ 
    and  $\bigwedge x. x \in \text{space } M \implies (\lambda i. s\ i\ x) \longrightarrow f\ x$ 
    and  $\bigwedge x\ i. x \in \text{space } M \implies \text{norm } (s\ i\ x) \leq 2 * \text{norm } (f\ x)$ 
proof -
  have  $f: f \in \text{borel-measurable } M (\int^+ x. \text{norm } (f\ x) \partial M) < \infty$  using assms
unfolding integrable-iff-bounded by auto
  obtain s where  $s: \bigwedge i. \text{simple-function } M (s\ i) \bigwedge x. x \in \text{space } M \implies (\lambda i. s\ i\ x) \longrightarrow f\ x$ 
     $\bigwedge i\ x. x \in \text{space } M \implies \text{norm } (s\ i\ x) \leq 2 * \text{norm } (f\ x)$  using
    borel-measurable-implies-sequence-metric[OF f(1)] unfolding norm-conv-dist by
    metis
  {
    fix i
    have  $(\int^+ x. \text{norm } (s\ i\ x) \partial M) \leq (\int^+ x. \text{ennreal } (2 * \text{norm } (f\ x)) \partial M)$  using
    s by (intro nn-integral-mono, auto)
    also have  $\dots < \infty$  using f by (simp add: nn-integral-cmult ennreal-mult-less-top ennreal-mult)
    finally have sbi: Bochner-Integration.simple-bochner-integrable M (s i) using
    s by (intro simple-bochner-integrableI-bounded) auto
    hence  $\text{emeasure } M \{y \in \text{space } M. s\ i\ y \neq 0\} \neq \infty$  by (auto intro: inte-

```

```

grableI-simple-bochner-integrable simple-bochner-integrable.cases)
}
thus ?thesis using that s by blast
qed

```

Simple functions can be represented by sums of indicator functions.

```

lemma simple-function-indicator-representation:
  fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
  assumes simple-function  $M f x \in \text{space } M$ 
  shows  $f x = (\sum y \in f \text{ ` space } M. \text{indicator } (f \text{ - ` } \{y\} \cap \text{space } M) x *_R y)$ 
  (is ?l = ?r)
proof -
  have ?r =  $(\sum y \in f \text{ ` space } M.$ 
     $(\text{if } y = f x \text{ then indicator } (f \text{ - ` } \{y\} \cap \text{space } M) x *_R y \text{ else } 0))$  by (auto intro!:
sum.cong)
  also have ... =  $\text{indicator } (f \text{ - ` } \{f x\} \cap \text{space } M) x *_R f x$  using assms by (auto
dest: simple-functionD)
  also have ... =  $f x$  using assms by (auto simp: indicator-def)
  finally show ?thesis by auto
qed

```

```

lemma simple-function-indicator-representation-AE:
  fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
  assumes simple-function  $M f$ 
  shows  $AE x \text{ in } M. f x = (\sum y \in f \text{ ` space } M. \text{indicator } (f \text{ - ` } \{y\} \cap \text{space } M) x$ 
 $*_R y)$ 
  by (metis (mono-tags, lifting) AE-I2 simple-function-indicator-representation
assms)

```

```

lemmas simple-function-scaleR[intro] = simple-function-compose2[where  $h = (*_R)$ ]

```

```

lemmas integrable-simple-function = simple-bochner-integrable.intros[THEN has-bochner-integral-simple-bochner]
THEN integrable.intros]

```

Induction rule for simple integrable functions.

```

lemma integrable-simple-function-induct[consumes 2, case-names cong indicator
add, induct set: simple-function]:
  fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
  assumes  $f: \text{simple-function } M f \text{ emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty$ 
  assumes  $\text{cong}: \bigwedge f g. \text{simple-function } M f \implies \text{emeasure } M \{y \in \text{space } M. f y \neq$ 
 $0\} \neq \infty$ 
 $\implies \text{simple-function } M g \implies \text{emeasure } M \{y \in \text{space } M. g y \neq$ 
 $0\} \neq \infty$ 
 $\implies (\bigwedge x. x \in \text{space } M \implies f x = g x) \implies P f \implies P g$ 
  assumes  $\text{indicator}: \bigwedge A y. A \in \text{sets } M \implies \text{emeasure } M A < \infty \implies P (\lambda x.$ 
 $\text{indicator } A x *_R y)$ 
  assumes  $\text{add}: \bigwedge f g. \text{simple-function } M f \implies \text{emeasure } M \{y \in \text{space } M. f y \neq$ 
 $0\} \neq \infty \implies$ 
 $\text{simple-function } M g \implies \text{emeasure } M \{y \in \text{space } M. g y \neq 0\} \neq$ 
 $\infty \implies$ 

```


$(\bigwedge z. z \in \text{space } M \implies \text{norm } (f z + g z) = \text{norm } (f z) + \text{norm } (g z)) \implies$
 $P f \implies P g \implies P (\lambda x. f x + g x)$
shows $P f$
proof –
let $?f = \lambda x. (\sum y \in f \text{ ` } \text{space } M. \text{indicat-real } (f - \text{ ` } \{y\} \cap \text{space } M) x *_R y)$
have $f\text{-ae-eq}: f x = ?f x$ **if** $x \in \text{space } M$ **for** x **using** $\text{simple-function-indicator-representation}[OF f(1) \text{ that}]$.
moreover have $\text{emeasure } M \{y \in \text{space } M. ?f y \neq 0\} \neq \infty$ **by** $(\text{metis } (\text{no-types}, \text{lifting}) \text{Collect-cong calculation } f(2))$
moreover have $P (\lambda x. \sum y \in S. \text{indicat-real } (f - \text{ ` } \{y\} \cap \text{space } M) x *_R y)$
 $\text{simple-function } M (\lambda x. \sum y \in S. \text{indicat-real } (f - \text{ ` } \{y\} \cap \text{space } M) x *_R y)$
 $\text{emeasure } M \{y \in \text{space } M. (\sum x \in S. \text{indicat-real } (f - \text{ ` } \{x\} \cap \text{space } M) y *_R x) \neq 0\} \neq \infty$
if $S \subseteq f \text{ ` } \text{space } M$ **for** S **using** $\text{simple-functionD}(1)[OF \text{assms}(1), \text{THEN rev-finite-subset, OF that}]$ **that**
proof $(\text{induction rule: finite-induct})$
case empty
 $\{$
case 1
then show $?case$ **using** $\text{indicator}[of \{ \}]$ **by force**
next
case 2
then show $?case$ **by force**
next
case 3
then show $?case$ **by force**
 $\}$
next
case $(\text{insert } x F)$
have $(f - \text{ ` } \{x\} \cap \text{space } M) \subseteq \{y \in \text{space } M. f y \neq 0\}$ **if** $x \neq 0$ **using** that **by blast**
moreover have $\{y \in \text{space } M. f y \neq 0\} = \text{space } M - (f - \text{ ` } \{0\} \cap \text{space } M)$
by blast
moreover have $\text{space } M - (f - \text{ ` } \{0\} \cap \text{space } M) \in \text{sets } M$ **using** $\text{simple-functionD}(2)[OF f(1)]$ **by blast**
ultimately have $\text{fin-0}: \text{emeasure } M (f - \text{ ` } \{x\} \cap \text{space } M) < \infty$ **if** $x \neq 0$ **using** that **by** $(\text{metis } \text{emeasure-mono } f(2) \text{infinity-ennreal-def top.not-eq-extremum top-unique})$
hence $\text{fin-1}: \text{emeasure } M \{y \in \text{space } M. \text{indicat-real } (f - \text{ ` } \{x\} \cap \text{space } M) y *_R x \neq 0\} \neq \infty$ **if** $x \neq 0$ **by** $(\text{metis } (\text{mono-tags}, \text{lifting}) \text{emeasure-mono } f(1) \text{indicator-simps}(2) \text{linorder-not-less mem-Collect-eq scaleR-eq-0-iff simple-functionD}(2) \text{subsetI that})$

have $*$: $(\sum y \in \text{insert } x F. \text{indicat-real } (f - \text{ ` } \{y\} \cap \text{space } M) x *_R y) = (\sum y \in F. \text{indicat-real } (f - \text{ ` } \{y\} \cap \text{space } M) x *_R y) + \text{indicat-real } (f - \text{ ` } \{x\} \cap \text{space } M) x *_R x$ **for** x **by** $(\text{metis } (\text{no-types}, \text{lifting}) \text{Diff-empty Diff-insert0 add commute insert.hyps}(1) \text{insert.hyps}(2) \text{sum.insert-remove})$

```

have **:  $\{y \in \text{space } M. (\sum x \in \text{insert } x \ F. \text{indicat-real } (f - \{x\} \cap \text{space } M) \ y$ 
 $*_R \ x) \neq 0\} \subseteq \{y \in \text{space } M. (\sum x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) \ y *_R \ x)$ 
 $\neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) \ y *_R \ x \neq 0\}$  unfolding
* by fastforce
{
  case 1
  hence  $x: x \in f \text{ ' space } M$  and  $F: F \subseteq f \text{ ' space } M$  by auto
  show ?case
  proof (cases  $x = 0$ )
    case True
    then show ?thesis unfolding * using insert(3)[OF F] by simp
  next
    case False
    have norm-argument:  $\text{norm } ((\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) \ z$ 
 $*_R \ y) + \text{indicat-real } (f - \{x\} \cap \text{space } M) \ z *_R \ x) = \text{norm } (\sum y \in F. \text{indicat-real}$ 
 $(f - \{y\} \cap \text{space } M) \ z *_R \ y) + \text{norm } (\text{indicat-real } (f - \{x\} \cap \text{space } M) \ z *_R \ x)$ 
if  $z: z \in \text{space } M$  for  $z$ 
    proof (cases  $f \ z = x$ )
      case True
      have  $\text{indicat-real } (f - \{y\} \cap \text{space } M) \ z *_R \ y = 0$  if  $y \in F$  for  $y$  using
True insert(2) z that 1 unfolding indicator-def by force
      hence  $(\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) \ z *_R \ y) = 0$  by (meson
sum.neutral)
      then show ?thesis by force
    next
      case False
      then show ?thesis by force
    qed
    show ?thesis using False simple-functionD(2)[OF f(1)] insert(3,5)[OF F]
simple-function-scaleR fin-0 fin-1 by (subst *, subst add, subst simple-function-sum)
(blast intro: norm-argument indicator)+
    qed
  next
    case 2
    hence  $x: x \in f \text{ ' space } M$  and  $F: F \subseteq f \text{ ' space } M$  by auto
    show ?case
    proof (cases  $x = 0$ )
      case True
      then show ?thesis unfolding * using insert(4)[OF F] by simp
    next
      case False
      then show ?thesis unfolding * using insert(4)[OF F] simple-functionD(2)[OF
f(1)] by fast
    qed
  next
    case 3
    hence  $x: x \in f \text{ ' space } M$  and  $F: F \subseteq f \text{ ' space } M$  by auto
    show ?case
    proof (cases  $x = 0$ )

```

```

case True
then show ?thesis unfolding * using insert(5)[OF F] by simp
next
case False
have emeasure M {y ∈ space M. (∑ x ∈ insert x F. indicat-real (f - ' {x} ∩ space M) y *R x) ≠ 0} ≤ emeasure M ({y ∈ space M. (∑ x ∈ F. indicat-real (f - ' {x} ∩ space M) y *R x) ≠ 0} ∪ {y ∈ space M. indicat-real (f - ' {x} ∩ space M) y *R x ≠ 0})
using ** simple-functionD(2)[OF insert(4)[OF F]] simple-functionD(2)[OF f(1)] by (intro emeasure-mono, force+)
also have ... ≤ emeasure M {y ∈ space M. (∑ x ∈ F. indicat-real (f - ' {x} ∩ space M) y *R x) ≠ 0} + emeasure M {y ∈ space M. indicat-real (f - ' {x} ∩ space M) y *R x ≠ 0}
using simple-functionD(2)[OF insert(4)[OF F]] simple-functionD(2)[OF f(1)] by (intro emeasure-subadditive, force+)
also have ... < ∞ using insert(5)[OF F] fin-1[OF False] by (simp add: less-top)
finally show ?thesis by simp
qed
}
qed
moreover have simple-function M (λx. ∑ y ∈ f ' space M. indicat-real (f - ' {y} ∩ space M) x *R y) using calculation by blast
moreover have P (λx. ∑ y ∈ f ' space M. indicat-real (f - ' {y} ∩ space M) x *R y) using calculation by blast
ultimately show ?thesis by (intro cong[OF - - f(1,2)], blast, presburger+)
qed

```

Induction rule for non-negative simple integrable functions

lemma *integrable-simple-function-induct-nn[consumes 3, case-names cong indicator add, induct set: simple-function]:*

fixes *f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, ordered-real-vector}*

assumes *f: simple-function M f emeasure M {y ∈ space M. f y ≠ 0} ≠ ∞ ∧ x. x ∈ space M ⇒ f x ≥ 0*

assumes *cong: ∧ f g. simple-function M f ⇒ emeasure M {y ∈ space M. f y ≠ 0} ≠ ∞ ⇒ (∧ x. x ∈ space M ⇒ f x ≥ 0) ⇒ simple-function M g ⇒ emeasure M {y ∈ space M. g y ≠ 0} ≠ ∞ ⇒ (∧ x. x ∈ space M ⇒ g x ≥ 0) ⇒ (∧ x. x ∈ space M ⇒ f x = g x) ⇒ P f ⇒ P g*

assumes *indicator: ∧ A y. y ≥ 0 ⇒ A ∈ sets M ⇒ emeasure M A < ∞ ⇒ P (λx. indicator A x *_R y)*

assumes *add: ∧ f g. (∧ x. x ∈ space M ⇒ f x ≥ 0) ⇒ simple-function M f ⇒ emeasure M {y ∈ space M. f y ≠ 0} ≠ ∞ ⇒*

(∧ x. x ∈ space M ⇒ g x ≥ 0) ⇒ simple-function M g ⇒ emeasure M {y ∈ space M. g y ≠ 0} ≠ ∞ ⇒

(∧ z. z ∈ space M ⇒ norm (f z + g z) = norm (f z) + norm (g z)) ⇒

P f ⇒ P g ⇒ P (λx. f x + g x)

shows *P f*

proof–

let $?f = \lambda x. (\sum_{y \in f^{-1} \text{ space } M} \text{indicat-real } (f - \{y\} \cap \text{space } M) \ x \ *_R y)$
have $f\text{-ae-eq}: f\ x = ?f\ x$ **if** $x \in \text{space } M$ **for** x **using** *simple-function-indicator-representation*[*OF* $f(1)$ *that*] .
moreover have $\text{emeasure } M \{y \in \text{space } M. ?f\ y \neq 0\} \neq \infty$ **by** (*metis* (*no-types*, *lifting*) *Collect-cong calculation* $f(2)$)
moreover have $P (\lambda x. \sum_{y \in S} \text{indicat-real } (f - \{y\} \cap \text{space } M) \ x \ *_R y)$
 $\text{simple-function } M (\lambda x. \sum_{y \in S} \text{indicat-real } (f - \{y\} \cap \text{space } M) \ x \ *_R y)$
 $\text{emeasure } M \{y \in \text{space } M. (\sum_{x \in S} \text{indicat-real } (f - \{x\} \cap \text{space } M) \ y \ *_R x) \neq 0\} \neq \infty$
 $\bigwedge x. x \in \text{space } M \implies 0 \leq (\sum_{y \in S} \text{indicat-real } (f - \{y\} \cap \text{space } M) \ x \ *_R y)$
if $S \subseteq f^{-1} \text{ space } M$ **for** S **using** *simple-functionD*(1)[*OF* *assms*(1), *THEN rev-finite-subset*, *OF* *that*] *that*
proof (*induction rule: finite-subset-induct'*)
case *empty*
{
case 1
then show $?case$ **using** *indicator*[*of* 0 { $\}$] **by** *force*
next
case 2
then show $?case$ **by** *force*
next
case 3
then show $?case$ **by** *force*
next
case 4
then show $?case$ **by** *force*
}
next
case (*insert* $x\ F$)
have $(f - \{x\} \cap \text{space } M) \subseteq \{y \in \text{space } M. f\ y \neq 0\}$ **if** $x \neq 0$ **using** *that* **by** *blast*
moreover have $\{y \in \text{space } M. f\ y \neq 0\} = \text{space } M - (f - \{0\} \cap \text{space } M)$
by *blast*
moreover have $\text{space } M - (f - \{0\} \cap \text{space } M) \in \text{sets } M$ **using** *simple-functionD*(2)[*OF* $f(1)$] **by** *blast*
ultimately have $\text{fin-0}: \text{emeasure } M (f - \{x\} \cap \text{space } M) < \infty$ **if** $x \neq 0$
using *that* **by** (*metis* *emeasure-mono* $f(2)$ *infinity-ennreal-def top.not-eq-extremum top-unique*)
hence $\text{fin-1}: \text{emeasure } M \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) \ y \ *_R x \neq 0\} \neq \infty$ **if** $x \neq 0$ **by** (*metis* (*mono-tags*, *lifting*) *emeasure-mono* $f(1)$ *indicator-simps*(2) *linorder-not-less mem-Collect-eq scaleR-eq-0-iff simple-functionD*(2) *subsetI* *that*)

have *nonneg-x*: $x \geq 0$ **using** *insert* f **by** *blast*

have $*$: $(\sum_{y \in \text{insert } x\ F} \text{indicat-real } (f - \{y\} \cap \text{space } M) \ x \ *_R y) = (\sum_{y \in F} \text{indicat-real } (f - \{y\} \cap \text{space } M) \ x \ *_R y) + \text{indicat-real } (f - \{x\} \cap \text{space } M) \ x \ *_R x$

$\text{space } M) \text{ } xa *_R x \text{ for } xa \text{ by } (\text{metis } (\text{no-types, lifting}) \text{ add.commute insert.hyps}(1) \text{ insert.hyps}(4) \text{ sum.insert})$
have **: $\{y \in \text{space } M. (\sum x \in \text{insert } x \text{ } F. \text{indicat-real } (f - \{x\} \cap \text{space } M) \text{ } y *_R x) \neq 0\} \subseteq \{y \in \text{space } M. (\sum x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) \text{ } y *_R x) \neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) \text{ } y *_R x \neq 0\}$ **unfolding**
*** by fastforce**
{
case 1
show ?case
proof (cases x = 0)
case True
then show ?thesis unfolding * using insert by simp
next
case False
have norm-argument: $\text{norm } ((\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) \text{ } z *_R y) + \text{indicat-real } (f - \{x\} \cap \text{space } M) \text{ } z *_R x) = \text{norm } (\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) \text{ } z *_R y) + \text{norm } (\text{indicat-real } (f - \{x\} \cap \text{space } M) \text{ } z *_R x)$
if z: $z \in \text{space } M$ **for z**
proof (cases f z = x)
case True
have $\text{indicat-real } (f - \{y\} \cap \text{space } M) \text{ } z *_R y = 0$ **if y ∈ F for y using**
True insert z that 1 **unfolding indicator-def by force**
hence $(\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) \text{ } z *_R y) = 0$ **by (meson**
sum.neutral)
thus ?thesis by force
qed (force)
show ?thesis using False fin-0 fin-1 f norm-argument by (subst *, subst add,
presburger add: insert, intro insert, intro insert, fastforce simp add: indicator-def
intro!: insert(2) f(3), auto intro!: indicator insert nonneg-x)
qed
next
case 2
show ?case
proof (cases x = 0)
case True
then show ?thesis unfolding * using insert by simp
next
case False
then show ?thesis unfolding * using insert simple-functionD(2)[OF f(1)]
by fast
qed
next
case 3
show ?case
proof (cases x = 0)
case True
then show ?thesis unfolding * using insert by simp
next
case False

```

      have emeasure M {y ∈ space M. (∑ x ∈ insert x F. indicat-real (f - ' {x}
      ∩ space M) y *R x) ≠ 0} ≤ emeasure M ({y ∈ space M. (∑ x ∈ F. indicat-real (f
      - ' {x} ∩ space M) y *R x) ≠ 0} ∪ {y ∈ space M. indicat-real (f - ' {x} ∩ space
      M) y *R x ≠ 0})
      using ** simple-functionD(2)[OF insert(6)] simple-functionD(2)[OF f(1)]
      insert.IH(2) by (intro emeasure-mono, blast, simp)
      also have ... ≤ emeasure M {y ∈ space M. (∑ x ∈ F. indicat-real (f - ' {x}
      ∩ space M) y *R x) ≠ 0} + emeasure M {y ∈ space M. indicat-real (f - ' {x} ∩
      space M) y *R x ≠ 0}
      using simple-functionD(2)[OF insert(6)] simple-functionD(2)[OF f(1)]
    by (intro emeasure-subadditive, force+)
      also have ... < ∞ using insert(7) fin-1[OF False] by (simp add: less-top)
      finally show ?thesis by simp
    qed
  next
    case (4 ξ)
    thus ?case using insert nonneg-x f(3) by (auto simp add: scaleR-nonneg-nonneg
    intro: sum-nonneg)
  }
  qed
  moreover have simple-function M (λx. ∑ y ∈ f ' space M. indicat-real (f - ' {y}
  ∩ space M) x *R y) using calculation by blast
  moreover have P (λx. ∑ y ∈ f ' space M. indicat-real (f - ' {y} ∩ space M) x
  *R y) using calculation by blast
  moreover have ∧x. x ∈ space M ⇒ 0 ≤ f x using f(3) by simp
  ultimately show ?thesis by (intro cong[OF - - - f(1,2)], blast, blast, fast)
presburger+
qed

lemma finite-nn-integral-imp-ae-finite:
  fixes f :: 'a ⇒ ennreal
  assumes f ∈ borel-measurable M (∫+x. f x ∂M) < ∞
  shows AE x in M. f x < ∞
proof (rule ccontr, goal-cases)
  case 1
  let ?A = space M ∩ {x. f x = ∞}
  have *: emeasure M ?A > 0 using 1 assms(1) by (metis (mono-tags, lifting)
  assms(2) eventually-mono infinity-ennreal-def nn-integral-noteq-infinite top.not-eq-extremum)
  have (∫+x. f x * indicator ?A x ∂M) = (∫+x. ∞ * indicator ?A x ∂M) by
  (metis (mono-tags, lifting) indicator-inter-arith indicator-simps(2) mem-Collect-eq
  mult-eq-0-iff)
  also have ... = ∞ * emeasure M ?A using assms(1) by (intro nn-integral-cmult-indicator,
  simp)
  also have ... = ∞ using * by fastforce
  finally have (∫+x. f x * indicator ?A x ∂M) = ∞ .
  moreover have (∫+x. f x * indicator ?A x ∂M) ≤ (∫+x. f x ∂M) by (intro
  nn-integral-mono, simp add: indicator-def)
  ultimately show ?case using assms(2) by simp
qed

```

Convergence in L1-Norm implies existence of a subsequence which convergences almost everywhere. This lemma is easier to use than the existing one in *HOL-Analysis.Bochner-Integration*

lemma *cauchy-L1-AE-cauchy-subseq*:

fixes $s :: \text{nat} \Rightarrow 'a \Rightarrow 'b :: \{\text{banach}, \text{second-countable-topology}\}$
assumes $[\text{measurable}] : \bigwedge n. \text{integrable } M (s \ n)$
and $\bigwedge e. e > 0 \implies \exists N. \forall i \geq N. \forall j \geq N. \text{LINT } x | M. \text{norm } (s \ i \ x - s \ j \ x) < e$
obtains r **where** *strict-mono* r *AE* x *in* M . *Cauchy* $(\lambda i. s \ (r \ i) \ x)$
proof –
have $\exists r. \forall n. (\forall i \geq r \ n. \forall j \geq r \ n. \text{LINT } x | M. \text{norm } (s \ i \ x - s \ j \ x) < (1 / 2) \wedge n) \wedge (r \ (\text{Suc } n) > r \ n)$
proof (*intro dependent-nat-choice, goal-cases*)
case 1
then show ?case **using** *assms*(2) **by** *presburger*
next
case (2 $x \ n$)
obtain N **where** *: $\text{LINT } x | M. \text{norm } (s \ i \ x - s \ j \ x) < (1 / 2) \wedge \text{Suc } n$ **if** $i \geq N \ j \geq N$ **for** $i \ j$ **using** *assms*(2)[*of* $(1 / 2) \wedge \text{Suc } n$] **by** *auto*
{
fix $i \ j$ **assume** $i \geq \max N \ (\text{Suc } x) \ j \geq \max N \ (\text{Suc } x)$
hence $\text{integral}^L M (\lambda x. \text{norm } (s \ i \ x - s \ j \ x)) < (1 / 2) \wedge \text{Suc } n$ **using** * **by** *fastforce*
}
then show ?case **by** *fastforce*
qed
then obtain r **where** *strict-mono*: *strict-mono* r **and** $\forall i \geq r \ n. \forall j \geq r \ n. \text{LINT } x | M. \text{norm } (s \ i \ x - s \ j \ x) < (1 / 2) \wedge n$ **for** n **using** *strict-mono-Suc-iff* **by** *blast*
hence *r-is*: $\text{LINT } x | M. \text{norm } (s \ (r \ (\text{Suc } n)) \ x - s \ (r \ n) \ x) < (1 / 2) \wedge n$ **for** n **by** (*simp add: strict-mono-leD*)

define g **where** $g = (\lambda n \ x. (\sum i \leq n. \text{ennreal } (\text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x))))$

define g' **where** $g' = (\lambda x. \sum i. \text{ennreal } (\text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x)))$

have *integrable-g*: $(\int^+ x. g \ n \ x \ \partial M) < 2$ **for** n

proof –

have $(\int^+ x. g \ n \ x \ \partial M) = (\int^+ x. (\sum i \leq n. \text{ennreal } (\text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x))) \ \partial M)$ **using** *g-def* **by** *simp*

also have ... = $(\sum i \leq n. (\int^+ x. \text{ennreal } (\text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x)) \ \partial M))$ **by** (*intro nn-integral-sum, simp*)

also have ... = $(\sum i \leq n. \text{LINT } x | M. \text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x))$ **unfolding** *dist-norm* **using** *assms*(1) **by** (*subst nn-integral-eq-integral[OF integrable-norm], auto*)

also have ... < $\text{ennreal } (\sum i \leq n. (1 / 2) \wedge i)$ **by** (*intro ennreal-lessI[OF sum-pos sum-strict-mono[OF finite-atMost - r-is]], auto*)

also have ... $\leq \text{ennreal } 2$ **unfolding** *sum-gp0*[*of* $1 / 2 \ n$] **by** (*intro ennreal-leI, auto*)

finally show $(\int^+ x. g \ n \ x \ \partial M) < 2$ **by** *simp*

qed

have *integrable-g'*: $(\int^+ x. g' x \partial M) \leq 2$
proof –
have *incseq* $(\lambda n. g n x)$ **for** x **by** (*intro incseq-SucI*, *auto simp add: g-def ennreal-leI*)
hence *convergent* $(\lambda n. g n x)$ **for** x **unfolding** *convergent-def* **using** *LIM-SEQ-incseq-SUP* **by** *fast*
hence $(\lambda n. g n x) \longrightarrow g' x$ **for** x **unfolding** *g-def g'-def* **by** (*intro summable-iff-convergent'[THEN iffD2, THEN summable-LIMSEQ]*, *blast*)
hence $(\int^+ x. g' x \partial M) = (\int^+ x. \liminf (\lambda n. g n x) \partial M)$ **by** (*metis lim-imp-Liminf trivial-limit-sequentially*)
also have $\dots \leq \liminf (\lambda n. \int^+ x. g n x \partial M)$ **by** (*intro nn-integral-liminf, simp add: g-def*)
also have $\dots \leq \liminf (\lambda n. 2)$ **using** *integrable-g* **by** (*intro Liminf-mono*) (*simp add: order-le-less*)
also have $\dots = 2$ **using** *sequentially-bot tendsto-iff-Liminf-eq-Limsup* **by** *blast*
finally show *?thesis* .
qed
hence *AE* x *in* $M. g' x < \infty$ **by** (*intro finite-nn-integral-imp-ae-finite*) (*auto simp add: order-le-less-trans g'-def*)
moreover have *summable* $(\lambda i. \text{norm } (s (r (Suc i)) x - s (r i) x))$ **if** $g' x \neq \infty$ **for** x **using** *that* **unfolding** *g'-def* **by** (*intro summable-suminf-not-top*) *fastforce* +

ultimately have *ae-summable*: *AE* x *in* $M. \text{summable } (\lambda i. s (r (Suc i)) x - s (r i) x)$ **using** *summable-norm-cancel* **unfolding** *dist-norm* **by** *force*

{
fix x **assume** *summable* $(\lambda i. s (r (Suc i)) x - s (r i) x)$
hence $(\lambda n. \sum_{i < n}. s (r (Suc i)) x - s (r i) x) \longrightarrow (\sum i. s (r (Suc i)) x - s (r i) x)$ **using** *summable-LIMSEQ* **by** *blast*
moreover have $(\lambda n. (\sum_{i < n}. s (r (Suc i)) x - s (r i) x)) = (\lambda n. s (r n) x - s (r 0) x)$ **using** *sum-lessThan-telescope* **by** *fastforce*
ultimately have $(\lambda n. s (r n) x - s (r 0) x) \longrightarrow (\sum i. s (r (Suc i)) x - s (r i) x)$ **by** *argo*
hence $(\lambda n. s (r n) x - s (r 0) x + s (r 0) x) \longrightarrow (\sum i. s (r (Suc i)) x - s (r i) x) + s (r 0) x$ **by** (*intro isCont-tendsto-compose[of - $\lambda z. z + s (r 0) x$], auto*)
hence *Cauchy* $(\lambda n. s (r n) x)$ **by** (*simp add: LIMSEQ-imp-Cauchy*)
}
hence *AE* x *in* $M. \text{Cauchy } (\lambda i. s (r i) x)$ **using** *ae-summable* **by** *fast*
thus *?thesis* **by** (*rule that[OF strict-mono(1)]*)
qed

3.2 Totally Ordered Banach Spaces

lemma *integrable-max*[*simp, intro*]:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology}\}$
assumes *fg*[*measurable*]: *integrable* $M f$ *integrable* $M g$
shows *integrable* $M (\lambda x. \max (f x) (g x))$
proof (*rule Bochner-Integration.integrable-bound*)


```

{
  fix x y :: 'b
  have norm (max x y) ≤ max (norm x) (norm y) by linarith
  also have ... ≤ norm x + norm y by simp
  finally have norm (max x y) ≤ norm (norm x + norm y) by simp
}
thus AE x in M. norm (max (f x) (g x)) ≤ norm (norm (f x) + norm (g x)) by
simp
qed (auto intro: Bochner-Integration.integrable-add[OF integrable-norm[OF fg(1)]
integrable-norm[OF fg(2)]])

```

```

lemma integrable-min[simp, intro]:
  fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology}
  assumes [measurable]: integrable M f integrable M g
  shows integrable M (λx. min (f x) (g x))
proof -
  have norm (min (f x) (g x)) ≤ norm (f x) ∨ norm (min (f x) (g x)) ≤ norm (g
x) for x by linarith
  thus ?thesis by (intro integrable-bound[OF integrable-max[OF integrable-norm(1,1),
OF assms]], auto)
qed

```

Restatement of *integral-nonneg-AE* for functions taking values in a Banach space.

```

lemma integral-nonneg-AE-banach:
  fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
dered-real-vector}
  assumes [measurable]: f ∈ borel-measurable M and nonneg: AE x in M. 0 ≤ f x
  shows 0 ≤ integralL M f
proof cases
  assume integrable: integrable M f
  hence max: (λx. max 0 (f x)) ∈ borel-measurable M ∧ x. 0 ≤ max 0 (f x)
integrable M (λx. max 0 (f x)) by auto
  hence 0 ≤ integralL M (λx. max 0 (f x))
  proof -
    obtain s where *: ∧i. simple-function M (s i)
      ∧i. emeasure M {y ∈ space M. s i y ≠ 0} ≠ ∞
      ∧x. x ∈ space M ⇒ (λi. s i x) ⟶ f x
      ∧x i. x ∈ space M ⇒ norm (s i x) ≤ 2 * norm (f x) using
integrable-implies-simple-function-sequence[OF integrable] by blast
    have simple: ∧i. simple-function M (λx. max 0 (s i x)) using * by fast
    have ∧i. {y ∈ space M. max 0 (s i y) ≠ 0} ⊆ {y ∈ space M. s i y ≠ 0}
  unfolding max-def by force
  moreover have ∧i. {y ∈ space M. s i y ≠ 0} ∈ sets M using * by measurable
  ultimately have ∧i. emeasure M {y ∈ space M. max 0 (s i y) ≠ 0} ≤
emeasure M {y ∈ space M. s i y ≠ 0} by (rule emeasure-mono)
  hence **: ∧i. emeasure M {y ∈ space M. max 0 (s i y) ≠ 0} ≠ ∞ using *(2)
by (auto intro: order.strict-trans1 simp add: top.not-eq-extremum)
  have ∧x. x ∈ space M ⇒ (λi. max 0 (s i x)) ⟶ max 0 (f x) using *(3)

```

tendsto-max **by** *blast*
moreover **have** $\bigwedge x i. x \in \text{space } M \implies \text{norm } (\max 0 (s i x)) \leq \text{norm } (2 *_R f x)$ **using** **(4) unfolding max-def by auto*
ultimately **have** *tendsto*: $(\lambda i. \text{integral}^L M (\lambda x. \max 0 (s i x))) \longrightarrow \text{integral}^L M (\lambda x. \max 0 (f x))$
using *borel-measurable-simple-function simple integrable by (intro integral-dominated-convergence[OF max(1) - integrable-norm[OF integrable-scaleR-right], of - 2 f], blast+)*
{
fix $h :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assume *simple-function* $M h \text{ emeasure } M \{y \in \text{space } M. h y \neq 0\} \neq \infty \bigwedge x. x \in \text{space } M \longrightarrow h x \geq 0$
hence $*: \text{integral}^L M h \geq 0$
proof (*induct rule: integrable-simple-function-induct-nn*)
case (*cong f g*)
then show *?case using Bochner-Integration.integral-cong by force*
next
case (*indicator A y*)
hence $A \neq \{\} \implies y \geq 0$ **using** *sets.sets-into-space by fastforce*
then show *?case using indicator by (cases A = {}, auto simp add: scaleR-nonneg-nonneg)*
next
case (*add f g*)
then show *?case by (simp add: integrable-simple-function)*
qed
}
thus *?thesis using LIMSEQ-le-const[OF tendsto, of 0] ** simple by fastforce*
qed
also **have** $\dots = \text{integral}^L M f$ **using** *nonneg by (auto intro: integral-cong-AE)*
finally **show** *?thesis .*
qed (*simp add: not-integrable-integral-eq*)

lemma *integral-mono-AE-banach*:

fixes $f g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes *integrable* $M f$ *integrable* $M g$ *AE* x *in* $M. f x \leq g x$
shows $\text{integral}^L M f \leq \text{integral}^L M g$
using *integral-nonneg-AE-banach[of $\lambda x. g x - f x$] assms Bochner-Integration.integral-diff[OF assms(1,2)] by force*

lemma *integral-mono-banach*:

fixes $f g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes *integrable* $M f$ *integrable* $M g$ $\bigwedge x. x \in \text{space } M \implies f x \leq g x$
shows $\text{integral}^L M f \leq \text{integral}^L M g$
using *integral-mono-AE-banach assms by blast*

3.3 Integrability and Measurability of the Diameter

context

fixes $s :: nat \Rightarrow 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$ **and** M

assumes $\text{bounded}: \bigwedge x. x \in \text{space } M \implies \text{bounded } (\text{range } (\lambda i. s \ i \ x))$

begin

lemma *borel-measurable-diameter*:

assumes $[\text{measurable}]: \bigwedge i. (s \ i) \in \text{borel-measurable } M$

shows $(\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\}) \in \text{borel-measurable } M$

proof –

have $\{s \ i \ x \mid i. N \leq i\} \neq \{\}$ **for** $x \ N$ **by** *blast*

hence *diameter-SUP*: $\text{diameter } \{s \ i \ x \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s \ i \ x) (s \ j \ x))$ **for** $x \ N$ **unfolding** *diameter-def* **by** $(\text{auto intro!}: \text{arg-cong}[\text{of} - \text{Sup}])$

have *case-prod dist* ‘ $(\{s \ i \ x \mid i. n \leq i\} \times \{s \ i \ x \mid i. n \leq i\}) = ((\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) ' (\{n..\} \times \{n..\}))$ ’ **for** x **by** *fast*

hence *: $(\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\}) = (\lambda x. \text{Sup } ((\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) ' (\{n..\} \times \{n..\})))$ **using** *diameter-SUP* **by** $(\text{simp add: case-prod-beta})$

have *bounded* $((\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) ' (\{n..\} \times \{n..\}))$ **if** $x \in \text{space } M$ **for** x **by** $(\text{rule bounded-imp-dist-bounded}[\text{OF bounded, OF that}])$

hence *bdd*: *bdd-above* $((\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) ' (\{n..\} \times \{n..\}))$ **if** $x \in \text{space } M$ **for** x **using** *that bounded-imp-bdd-above* **by** *presburger*

have *fst* $p \in \text{borel-measurable } M$ **and** $p \in \text{borel-measurable } M$ **if** $p \in s ' \{n..\} \times s ' \{n..\}$ **for** p **using** *that* **by** *fastforce+*

hence $(\lambda x. \text{fst } p \ x - \text{snd } p \ x) \in \text{borel-measurable } M$ **if** $p \in s ' \{n..\} \times s ' \{n..\}$ **for** p **using** *that borel-measurable-diff* **by** *simp*

hence $(\lambda x. \text{case } p \text{ of } (f, g) \Rightarrow \text{dist } (f \ x) (g \ x)) \in \text{borel-measurable } M$ **if** $p \in s ' \{n..\} \times s ' \{n..\}$ **for** p **unfolding** *dist-norm* **using** *that* **by** *measurable*

moreover **have** *countable* $(s ' \{n..\} \times s ' \{n..\})$ **by** $(\text{intro countable-SIGMA countable-image, auto})$

ultimately show *?thesis* **unfolding** * **by** $(\text{auto intro!}: \text{borel-measurable-cSUP bdd})$

qed

lemma *integrable-bound-diameter*:

fixes $f :: 'a \Rightarrow \text{real}$

assumes *integrable* $M \ f$

and $[\text{measurable}]: \bigwedge i. (s \ i) \in \text{borel-measurable } M$

and $\bigwedge x \ i. x \in \text{space } M \implies \text{norm } (s \ i \ x) \leq f \ x$

shows *integrable* $M \ (\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\})$

proof –

have $\{s \ i \ x \mid i. N \leq i\} \neq \{\}$ **for** $x \ N$ **by** *blast*

hence *diameter-SUP*: $\text{diameter } \{s \ i \ x \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s \ i \ x) (s \ j \ x))$ **for** $x \ N$ **unfolding** *diameter-def* **by** $(\text{auto intro!}: \text{arg-cong}[\text{of} - \text{Sup}])$

{

fix x **assume** $x: x \in \text{space } M$

```

    let ?S = (λ(i, j). dist (s i x) (s j x)) ‘ ({n..} × {n..})
    have case-prod dist ‘ ({s i x | i. n ≤ i} × {s i x | i. n ≤ i}) = (λ(i, j). dist (s i
x) (s j x)) ‘ ({n..} × {n..}) by fast
    hence *: diameter {s i x | i. n ≤ i} = Sup ?S using diameter-SUP by (simp
add: case-prod-beta')

    have bounded ?S by (rule bounded-imp-dist-bounded[OF bounded[OF x]])
    hence Sup-S-nonneg: 0 ≤ Sup ?S by (auto intro!: cSup-upper2 x bounded-imp-bdd-above)

    have dist (s i x) (s j x) ≤ 2 * f x for i j by (intro dist-triangle2[THEN
order-trans, of - 0]) (metis norm-conv-dist assms(3) x add-mono mult-2)
    hence ∀ c ∈ ?S. c ≤ 2 * f x by force
    hence Sup ?S ≤ 2 * f x by (intro cSup-least, auto)
    hence norm (Sup ?S) ≤ 2 * norm (f x) using Sup-S-nonneg by auto
    also have ... = norm (2 *R f x) by simp
    finally have norm (diameter {s i x | i. n ≤ i}) ≤ norm (2 *R f x) unfolding
* .
}
    hence AE x in M. norm (diameter {s i x | i. n ≤ i}) ≤ norm (2 *R f x) by blast
    thus integrable M (λx. diameter {s i x | i. n ≤ i}) using borel-measurable-diameter
by (intro Bochner-Integration.integrable-bound[OF assms(1)[THEN integrable-scaleR-right[of
2]]], measurable)
qed
end

```

3.4 Auxiliary Lemmas for Set Integrals

lemma *set-integral-scaleR-left*:

```

assumes A ∈ sets M c ≠ 0 ⇒ integrable M f
shows LINT t:A|M. f t *R c = (LINT t:A|M. f t) *R c
unfolding set-lebesgue-integral-def
using integrable-mult-indicator[OF assms]
by (subst integral-scaleR-left[symmetric], auto)

```

lemma *nn-set-integral-eq-set-integral*:

```

assumes [measurable]: integrable M f
and AE x ∈ A in M. 0 ≤ f x A ∈ sets M
shows (∫+ x ∈ A. f x ∂M) = (∫ x ∈ A. f x ∂M)

```

proof –

```

    have (∫+ x. indicator A x *R f x ∂M) = (∫ x ∈ A. f x ∂M)
    unfolding set-lebesgue-integral-def using assms(2) by (intro nn-integral-eq-integral[of
- λx. indicat-real A x *R f x], blast intro: assms integrable-mult-indicator, fastforce)
    moreover have (∫+ x. indicator A x *R f x ∂M) = (∫+ x ∈ A. f x ∂M) by (metis
ennreal-0 indicator-simps(1) indicator-simps(2) mult.commute mult-1 mult-zero-left
real-scaleR-def)

```

ultimately show ?thesis by argo

qed

lemma *set-integral-restrict-space*:

```

fixes f :: 'a  $\Rightarrow$  'b::{banach, second-countable-topology}
assumes  $\Omega \cap \text{space } M \in \text{sets } M$ 
shows set-lebesgue-integral (restrict-space M  $\Omega$ ) A f = set-lebesgue-integral M A
( $\lambda x. \text{indicator } \Omega \ x *_{\mathbb{R}} f \ x$ )
unfolding set-lebesgue-integral-def
by (subst integral-restrict-space, auto intro!: integrable-mult-indicator assms simp:
mult.commute)

lemma set-integral-const:
fixes c :: 'b::{banach, second-countable-topology}
assumes A  $\in \text{sets } M$  emeasure M A  $\neq \infty$ 
shows set-lebesgue-integral M A ( $\lambda \cdot. c$ ) = measure M A  $*_{\mathbb{R}}$  c
unfolding set-lebesgue-integral-def
using assms by (metis has-bochner-integral-indicator has-bochner-integral-integral-eq
infinity-enreal-def less-top)

lemma set-integral-mono-banach:
fixes f g :: 'a  $\Rightarrow$  'b :: {second-countable-topology, banach, linorder-topology, or-
dered-real-vector}
assumes set-integrable M A f set-integrable M A g
 $\bigwedge x. x \in A \implies f \ x \leq g \ x$ 
shows (LINT x:A|M. f x)  $\leq$  (LINT x:A|M. g x)
using assms unfolding set-integrable-def set-lebesgue-integral-def
by (auto intro: integral-mono-banach split: split-indicator)

lemma set-integral-mono-AE-banach:
fixes f g :: 'a  $\Rightarrow$  'b :: {second-countable-topology, banach, linorder-topology, or-
dered-real-vector}
assumes set-integrable M A f set-integrable M A g AE x $\in$ A in M. f x  $\leq$  g x
shows set-lebesgue-integral M A f  $\leq$  set-lebesgue-integral M A g using assms
unfolding set-lebesgue-integral-def by (auto simp add: set-integrable-def intro!:
integral-mono-AE-banach[of M  $\lambda x. \text{indicator } A \ x *_{\mathbb{R}} f \ x \ \lambda x. \text{indicator } A \ x *_{\mathbb{R}} g \ x$ ],
simp add: indicator-def)

```

3.5 Averaging Theorem

We aim to lift results from the real case to arbitrary Banach spaces. Our fundamental tool in this regard will be the averaging theorem. The proof of this theorem is due to Serge Lang (Real and Functional Analysis) \cite{Lang-1993}. The theorem allows us to make statements about a functions value almost everywhere, depending on the value its integral takes on various sets of the measure space.

Before we introduce and prove the averaging theorem, we will first show the following lemma which is crucial for our proof. While not stated exactly in this manner, our proof makes use of the characterization of second-countable topological spaces given in the book General Topology by Ryszard Engelking (Theorem 4.1.15) \cite{engelking-1989}.

lemma *balls-countable-basis*:

obtains $D :: 'a :: \{\text{metric-space, second-countable-topology}\}$ *set*
where *topological-basis* (*case-prod* ball ' ($D \times (\mathbb{Q} \cap \{0 < ..\})$)))
and *countable* D
and $D \neq \{\}$
proof –
obtain $D :: 'a$ *set* **where** *dense-subset*: *countable* D $D \neq \{\}$ $\llbracket \text{open } U; U \neq \{\} \rrbracket$
 $\implies \exists y \in D. y \in U$ **for** U **using** *countable-dense-exists* **by** *blast*
have *topological-basis* (*case-prod* ball ' ($D \times (\mathbb{Q} \cap \{0 < ..\})$)))
proof (*intro* *topological-basis-iff*[*THEN* *iffD2*], *fast*, *clarify*)
fix U **and** $x :: 'a$ **assume** *asm*: *open* U $x \in U$
obtain e **where** $e: e > 0$ *ball* $x \in U$ **using** *asm* *openE* **by** *blast*
obtain y **where** $y: y \in D$ $y \in \text{ball } x (e / 3)$ **using** *dense-subset*(3)[*OF* *open-ball*,
of $x \in U$ / 3] *centre-in-ball*[*THEN* *iffD2*, *OF* *divide-pos-pos*[*OF* $e(1)$, *of* 3]] **by** *force*
obtain r **where** $r: r \in \mathbb{Q} \cap \{e/3 < .. < e/2\}$ **unfolding** *Rats-def* **using** *of-rat-dense*[*OF*
divide-strict-left-mono[*OF* - $e(1)$], *of* 2 3] **by** *auto*

have $*$: $x \in \text{ball } y r$ **using** r y **by** (*simp* *add*: *dist-commute*)
hence $\text{ball } y r \subseteq U$ **using** r **by** (*intro* *order-trans*[*OF* - $e(2)$], *simp*, *metric*)
moreover **have** $\text{ball } y r \in (\text{case-prod } \text{ball ' } (D \times (\mathbb{Q} \cap \{0 < ..\})))$ **using** $y(1)$
r **by** *force*
ultimately show $\exists B' \in (\text{case-prod } \text{ball ' } (D \times (\mathbb{Q} \cap \{0 < ..\}))). x \in B' \wedge B' \subseteq$
 U **using** $*$ **by** *meson*
qed
thus *?thesis* **using** *that* *dense-subset* **by** *blast*
qed

context *sigma-finite-measure*
begin

To show statements concerning σ -finite measure spaces, one usually shows the statement for finite measure spaces and uses a limiting argument to show it for the σ -finite case. The following induction scheme formalizes this.

lemma *sigma-finite-measure-induct*[*case-names* *finite-measure*, *consumes* 0]:

assumes $\bigwedge(N :: 'a \text{ measure}) \Omega. \text{finite-measure } N$
 $\implies N = \text{restrict-space } M \Omega$
 $\implies \Omega \in \text{sets } M$
 $\implies \text{emeasure } N \Omega \neq \infty$
 $\implies \text{emeasure } N \Omega \neq 0$
 $\implies \text{almost-everywhere } N Q$
and [*measurable*]: *Measurable.pred* $M Q$
shows *almost-everywhere* $M Q$
proof –
have $*$: *almost-everywhere* $N Q$ **if** *finite-measure* N $N = \text{restrict-space } M \Omega$ $\Omega \in \text{sets } M$ $\text{emeasure } N \Omega \neq \infty$ **for** $N \Omega$ **using** *that* **by** (*cases* $\text{emeasure } N \Omega = 0$,
auto *intro*: *emeasure-0-AE* *assms*(1))

obtain $A :: \text{nat} \Rightarrow 'a$ *set* **where** A : *range* $A \subseteq \text{sets } M$ $(\bigcup i. A i) = \text{space } M$ **and**
emeasure-finite: $\text{emeasure } M (A i) \neq \infty$ **for** i **using** *sigma-finite* **by** *metis*

```

note  $A(1)[\text{measurable}]$ 
have  $\text{space-restr} : \text{space } (\text{restrict-space } M (A i)) = A i$  for  $i$  unfolding  $\text{space-restrict-space}$ 
by  $\text{simp}$ 
{
  fix  $i$ 
  have  $*$ :  $\{x \in A i \cap \text{space } M. Q x\} = \{x \in \text{space } M. Q x\} \cap (A i)$  by  $\text{fast}$ 
  have  $\text{Measurable.pred } (\text{restrict-space } M (A i)) Q$  using  $A$  by  $(\text{intro measurableI},$ 
 $\text{auto simp add: space-restr intro!: sets-restrict-space-iff[THEN iffD2], measurable},$ 
 $\text{auto})$ 
}
note  $\text{this}[\text{measurable}]$ 
{
  fix  $i$ 
  have  $\text{finite-measure } (\text{restrict-space } M (A i))$  using  $\text{emeasure-finite}$  by  $(\text{intro}$ 
 $\text{finite-measureI}, \text{subst space-restr}, \text{subst emeasure-restrict-space}, \text{auto})$ 
  hence  $\text{emeasure } (\text{restrict-space } M (A i)) \{x \in A i. \neg Q x\} = 0$  using  $\text{emea-}$ 
 $\text{sure-finite}$  by  $(\text{intro AE-iff-measurable[THEN iffD1, OF - - *]}, \text{measurable}, \text{subst}$ 
 $\text{space-restr[symmetric]}, \text{intro sets.top}, \text{auto simp add: emeasure-restrict-space})$ 
  hence  $\text{emeasure } M \{x \in A i. \neg Q x\} = 0$  by  $(\text{subst emeasure-restrict-space[symmetric]},$ 
 $\text{auto})$ 
}
hence  $\text{emeasure } M (\bigcup i. \{x \in A i. \neg Q x\}) = 0$  by  $(\text{intro emeasure-UN-eq-0},$ 
 $\text{auto})$ 
moreover have  $(\bigcup i. \{x \in A i. \neg Q x\}) = \{x \in \text{space } M. \neg Q x\}$  using  $A$  by
 $\text{auto}$ 
ultimately show  $?thesis$  by  $(\text{intro AE-iff-measurable[THEN iffD2]}, \text{auto})$ 
qed

```

The Averaging Theorem allows us to make statements concerning how a function behaves almost everywhere, depending on its behaviour on average.

lemma *averaging-theorem*:

```

fixes  $f :: \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
assumes  $[\text{measurable}]: \text{integrable } M f$ 
and  $\text{closed}: \text{closed } S$ 
and  $\bigwedge A. A \in \text{sets } M \implies \text{measure } M A > 0 \implies (1 / \text{measure } M A) *_R$ 
 $\text{set-lebesgue-integral } M A f \in S$ 
shows  $\text{AE } x \text{ in } M. f x \in S$ 
proof  $(\text{induct rule: sigma-finite-measure-induct})$ 
case  $(\text{finite-measure } N \Omega)$ 

```

interpret $\text{finite-measure } N$ **by** $(\text{rule finite-measure})$

```

have  $\text{integrable}[\text{measurable}]: \text{integrable } N f$  using  $\text{assms finite-measure}$  by  $(\text{auto}$ 
 $\text{simp: integrable-restrict-space integrable-mult-indicator})$ 
have  $\text{average}: (1 / \text{Sigma-Algebra.measure } N A) *_R \text{set-lebesgue-integral } N A f$ 
 $\in S$  if  $A \in \text{sets } N \text{ measure } N A > 0$  for  $A$ 
proof  $-$ 
have  $*$ :  $A \in \text{sets } M$  using  $\text{that}$  by  $(\text{simp add: sets-restrict-space-iff finite-measure})$ 
have  $A = A \cap \Omega$  by  $(\text{metis finite-measure}(2,3) \text{inf.orderE sets.sets-into-space})$ 

```

space-restrict-space that(1))

hence $\text{set-lebesgue-integral } N \ A \ f = \text{set-lebesgue-integral } M \ A \ f$ **unfolding**
finite-measure **by** (*subst set-integral-restrict-space*, *auto simp add: finite-measure*
set-lebesgue-integral-def indicator-inter-arith[symmetric])

moreover **have** $\text{measure } N \ A = \text{measure } M \ A$ **using** *that* **by** (*auto intro!:*
measure-restrict-space simp add: finite-measure sets-restrict-space-iff)

ultimately **show** *?thesis* **using** *that * assms(3)* **by** *presburger*

qed

obtain $D :: 'b \text{ set}$ **where** *balls-basis: topological-basis (case-prod ball ' (D × (Q*
 $\cap \{0 < ..\}))$) **and** *countable-D: countable D* **using** *balls-countable-basis* **by** *blast*

have *countable-balls: countable (case-prod ball ' (D × (Q*
 $\cap \{0 < ..\}))$) **using**
countable-rat countable-D **by** *blast*

obtain B **where** *B-balls: B ⊆ case-prod ball ' (D × (Q*
 $\cap \{0 < ..\})) \cup B = -S$
using *topological-basis[THEN iffD1, OF balls-basis] open-Compl[OF assms(2)]* **by**
meson

hence *countable-B: countable B* **using** *countable-balls countable-subset* **by** *fast*

define b **where** $b = \text{from-nat-into } (B \cup \{\{\}\})$

have $B \cup \{\{\}\} \neq \{\}$ **by** *simp*

have *range-b: range b = B ∪ { { } }* **using** *countable-B* **by** (*auto simp add: b-def*
intro!: *range-from-nat-into*)

have *open-b: open (b i) for i* **unfolding** *b-def* **using** *B-balls open-ball from-nat-into[of*
 $B \cup \{\{\}\}$ $i]$ **by** *force*

have *Union-range-b: ∪ (range b) = -S* **using** *B-balls range-b* **by** *simp*

{

fix $v \ r$ **assume** *ball-in-Compl: ball v r ⊆ -S*

define A **where** $A = f - ' \text{ball } v \ r \cap \text{space } N$

have *dist-less: dist (f x) v < r if x ∈ A for x* **using** *that* **unfolding** *A-def*
vimage-def **by** (*simp add: dist-commute*)

hence *AE-less: AE x ∈ A in N. norm (f x - v) < r* **by** (*auto simp add:*
dist-norm)

have $*$: $A \in \text{sets } N$ **unfolding** *A-def* **by** *simp*

have $\text{emeasure } N \ A = 0$

proof -

{

assume *asm: emeasure N A > 0*

hence *measure-pos: measure N A > 0* **unfolding** *emeasure-eq-measure* **by**
simp

hence $(1 / \text{measure } N \ A) *_{\mathbb{R}} \text{set-lebesgue-integral } N \ A \ f - v = (1 / \text{measure } N$
 $A) *_{\mathbb{R}} \text{set-lebesgue-integral } N \ A \ (\lambda x. f \ x - v)$ **using** *integrable integrable-const ** **by**
*(subst set-integral-diff(2), auto simp add: set-integrable-def set-integral-const[OF *]*
algebra-simps intro!: *integrable-mult-indicator*)

moreover **have** $\text{norm } (\int x \in A. (f \ x - v) \partial N) \leq (\int x \in A. \text{norm } (f \ x$
 $- v) \partial N)$ **using** $*$ **by** (*auto intro!:* *integral-norm-bound[of N λx. indicator A x*
 $*_{\mathbb{R}} (f \ x - v), \text{ THEN order-trans}]$ *integrable-mult-indicator integrable simp add:*
set-lebesgue-integral-def)

ultimately have $\text{norm } ((1 / \text{measure } N \ A) *_{\mathcal{R}} \text{set-lebesgue-integral } N \ A \ f - v) \leq \text{set-lebesgue-integral } N \ A \ (\lambda x. \text{norm } (f \ x - v)) / \text{measure } N \ A$ **using** *asm*
by (*auto intro: divide-right-mono*)
also have $\dots < \text{set-lebesgue-integral } N \ A \ (\lambda x. r) / \text{measure } N \ A$
unfolding *set-lebesgue-integral-def*
using *asm * integrable integrable-const AE-less measure-pos*
by (*intro divide-strict-right-mono integral-less-AE[of - - A] integrable-mult-indicator*)
(fastforce simp add: dist-less dist-norm indicator-def)+
also have $\dots = r$ **using** ** measure-pos* **by** (*simp add: set-integral-const*)
finally have $\text{dist } ((1 / \text{measure } N \ A) *_{\mathcal{R}} \text{set-lebesgue-integral } N \ A \ f) \ v < r$
by (*subst dist-norm*)
hence *False* **using** *average[OF * measure-pos]* **by** (*metis ComplD dist-commute in-mono mem-ball ball-in-Compl*)
}
thus *?thesis* **by** *fastforce*
qed
}
note $*$ **=** *this*
{
fix b' **assume** $b' \in B$
hence *ball-subset-Compl*: $b' \subseteq -S$ **and** *ball-radius-pos*: $\exists v \in D. \exists r > 0. b' = \text{ball } v \ r$ **using** *B-balls* **by** (*blast, fast*)
}
note $**$ **=** *this*
hence $\text{emeasure } N \ (f - ' b \ i \cap \text{space } N) = 0$ **for** i **by** (*cases b i = {}, simp*)
*(metis UnE singletonD * range-b[THEN eq-refl, THEN range-subsetD])*
hence $\text{emeasure } N \ (\bigcup i. f - ' b \ i \cap \text{space } N) = 0$ **using** *open-b* **by** (*intro emeasure-UN-eq-0 fastforce+*)
moreover have $(\bigcup i. f - ' b \ i \cap \text{space } N) = f - ' (\bigcup (\text{range } b)) \cap \text{space } N$ **by** *blast*
ultimately have $\text{emeasure } N \ (f - ' (-S) \cap \text{space } N) = 0$ **using** *Union-range-b*
by *argo*
hence $\text{AE } x \text{ in } N. f \ x \notin -S$ **using** *open-Compl[OF assms(2)]* **by** (*intro AE-iff-measurable[THEN iffD2], auto*)
thus *?case* **by** *force*
qed (*simp add: pred-sets2[OF borel-closed] assms(2)*)

lemma *density-zero*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes *integrable M f*
and *density-0*: $\bigwedge A. A \in \text{sets } M \implies \text{set-lebesgue-integral } M \ A \ f = 0$
shows $\text{AE } x \text{ in } M. f \ x = 0$
using *averaging-theorem[OF assms(1), of {0}] assms(2)*
by (*simp add: scaleR-nonneg-nonneg*)

The following lemma shows that densities are unique in Banach spaces.

lemma *density-unique-banach*:

fixes $f \ f' :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes *integrable M f integrable M f'*

and *density-eq*: $\bigwedge A. A \in \text{sets } M \implies \text{set-lebesgue-integral } M A f = \text{set-lebesgue-integral } M A f'$
shows $AE\ x\ \text{in } M. f\ x = f'\ x$
proof –
 {
 fix A **assume** $asm: A \in \text{sets } M$
 hence $LINT\ x|M. \text{indicat-real } A\ x *_R (f\ x - f'\ x) = 0$ **using** *density-eq*
 $assms(1,2)$ **by** (*simp add: set-lebesgue-integral-def algebra-simps Bochner-Integration.integral-diff[OF integrable-mult-indicator(1,1)]*)
 }
thus *?thesis* **using** *density-zero[OF Bochner-Integration.integrable-diff[OF assms(1,2)]]*
by (*simp add: set-lebesgue-integral-def*)
qed

lemma *density-nonneg*:
fixes $f :: - \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes *integrable* $M f$
and $\bigwedge A. A \in \text{sets } M \implies \text{set-lebesgue-integral } M A f \geq 0$
shows $AE\ x\ \text{in } M. f\ x \geq 0$
using *averaging-theorem[OF assms(1), of \{0..\}, OF closed-atLeast] assms(2)*
by (*simp add: scaleR-nonneg-nonneg*)

corollary *integral-nonneg-eq-0-iff-AE-banach*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes $f[\text{measurable}]: \text{integrable } M f$ **and** *nonneg*: $AE\ x\ \text{in } M. 0 \leq f\ x$
shows $\text{integral}^L M f = 0 \longleftrightarrow (AE\ x\ \text{in } M. f\ x = 0)$

proof

assume $*$: $\text{integral}^L M f = 0$
 {
 fix A **assume** $asm: A \in \text{sets } M$
 have $0 \leq \text{integral}^L M (\lambda x. \text{indicator } A\ x *_R f\ x)$ **using** *nonneg* **by** (*subst integral-zero[of M, symmetric], intro integral-mono-AE-banach integrable-mult-indicator asm f integrable-zero, auto simp add: indicator-def*)
 moreover **have** $\dots \leq \text{integral}^L M f$ **using** *nonneg* **by** (*intro integral-mono-AE-banach integrable-mult-indicator asm f, auto simp add: indicator-def*)
 ultimately **have** $\text{set-lebesgue-integral } M A f = 0$ **unfolding** *set-lebesgue-integral-def*
 using $*$ **by** *force*
 }
thus $AE\ x\ \text{in } M. f\ x = 0$ **by** (*intro density-zero f, blast*)
qed (*auto simp add: integral-eq-zero-AE*)

corollary *integral-eq-mono-AE-eq-AE*:

fixes $f\ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes *integrable* $M f$ *integrable* $M g$ $\text{integral}^L M f = \text{integral}^L M g$ $AE\ x\ \text{in } M. f\ x \leq g\ x$
shows $AE\ x\ \text{in } M. f\ x = g\ x$
proof –

```

define  $h$  where  $h = (\lambda x. g\ x - f\ x)$ 
  have  $AE\ x\ in\ M. h\ x = 0$  unfolding  $h\text{-def}$  using  $assms$  by ( $subst\ integr\text{-}nonneg\text{-}eq\text{-}0\text{-}iff\text{-}AE\text{-}banach[symmetric]$ )  $auto$ 
  then show  $?thesis$  unfolding  $h\text{-def}$  by  $auto$ 
qed

end

end

```

```

theory Conditional-Expectation-Banach
  imports HOL-Probability.Conditional-Expectation HOL-Probability.Independent-Family
  Bochner-Integration-Supplement
begin

```

4 Conditional Expectation in Banach Spaces

While constructing the conditional expectation operator, we have come up with the following approach, which is based on the construction in [1]. Both our approach, and the one in [1] are based on showing that the conditional expectation is a contraction on some dense subspace of the space of functions $L^1(E)$. In our approach, we start by constructing the conditional expectation explicitly for simple functions. Then we show that the conditional expectation is a contraction on simple functions, i.e. $\|E(s|F)(x)\| \leq E(\|s(x)\||F)$ for μ -almost all $x \in \Omega$ with $s : \Omega \rightarrow E$ simple and integrable. Using this, we can show that the conditional expectation of a convergent sequence of simple functions is again convergent. Finally, we show that this limit exhibits the properties of a conditional expectation. This approach has the benefit of being straightforward and easy to implement, since we could make use of the existing formalization for real-valued functions. To use the construction in [1] we need more tools from functional analysis, which Isabelle/HOL currently does not have.

Before we can talk about 'the' conditional expectation, we must define what it means for a function to have a conditional expectation.

definition $has\text{-}cond\text{-}exp :: 'a\ measure \Rightarrow 'a\ measure \Rightarrow ('a \Rightarrow 'b) \Rightarrow ('a \Rightarrow 'b :: \{real\text{-}normed\text{-}vector, second\text{-}countable\text{-}topology\}) \Rightarrow bool$ **where**
 $has\text{-}cond\text{-}exp\ M\ F\ f\ g = ((\forall A \in sets\ F. (\int x \in A. f\ x\ \partial M) = (\int x \in A. g\ x\ \partial M))$

$$\begin{aligned}
&\wedge integrable\ M\ f \\
&\wedge integrable\ M\ g \\
&\wedge g \in borel\text{-}measurable\ F)
\end{aligned}$$

This predicate precisely characterizes what it means for a function f to have a conditional expectation g , with respect to the measure M and the sub-

algebra F .

lemma *has-cond-expI'*:

assumes $\bigwedge A. A \in \text{sets } F \implies (\int x \in A. f x \partial M) = (\int x \in A. g x \partial M)$
 $\text{integrable } M f$
 $\text{integrable } M g$
 $g \in \text{borel-measurable } F$
shows $\text{has-cond-exp } M F f g$
using *assms unfolding has-cond-exp-def by simp*

lemma *has-cond-expD*:

assumes $\text{has-cond-exp } M F f g$
shows $\bigwedge A. A \in \text{sets } F \implies (\int x \in A. f x \partial M) = (\int x \in A. g x \partial M)$
 $\text{integrable } M f$
 $\text{integrable } M g$
 $g \in \text{borel-measurable } F$
using *assms unfolding has-cond-exp-def by simp+*

Now we can use Hilberts ϵ -operator to define the conditional expectation, if it exists.

definition $\text{cond-exp} :: 'a \text{ measure} \Rightarrow 'a \text{ measure} \Rightarrow ('a \Rightarrow 'b) \Rightarrow ('a \Rightarrow 'b :: \{\text{banach, second-countable-topology}\})$ **where**

$\text{cond-exp } M F f = (\text{if } \exists g. \text{has-cond-exp } M F f g \text{ then } (\text{SOME } g. \text{has-cond-exp } M F f g) \text{ else } (\lambda -. 0))$

lemma *borel-measurable-cond-exp[measurable]*: $\text{cond-exp } M F f \in \text{borel-measurable } F$

by (*metis cond-exp-def someI has-cond-exp-def borel-measurable-const*)

lemma *integrable-cond-exp[intro]*: $\text{integrable } M (\text{cond-exp } M F f)$

by (*metis cond-exp-def has-cond-expD(3) integrable-zero someI*)

lemma *set-integrable-cond-exp[intro]*:

assumes $A \in \text{sets } M$

shows $\text{set-integrable } M A (\text{cond-exp } M F f)$ **using** *integrable-mult-indicator[OF assms integrable-cond-exp, of F f]* **by** (*auto simp add: set-integrable-def intro!: integrable-mult-indicator[OF assms integrable-cond-exp]*)

lemma *has-cond-exp-self*:

assumes $\text{integrable } M f$

shows $\text{has-cond-exp } M (\text{vimage-algebra } (\text{space } M) f \text{ borel}) f f$

using *assms by (auto intro!: has-cond-expI' measurable-vimage-algebra1)*

lemma *has-cond-exp-sets-cong*:

assumes $\text{sets } F = \text{sets } G$

shows $\text{has-cond-exp } M F = \text{has-cond-exp } M G$

using *assms unfolding has-cond-exp-def by force*

lemma *cond-exp-sets-cong*:

assumes $\text{sets } F = \text{sets } G$

shows $AE\ x\ in\ M.\ cond_exp\ M\ F\ f\ x = cond_exp\ M\ G\ f\ x$
by (intro AE-I2, simp add: cond-exp-def has-cond-exp-sets-cong[OF assms, of M])

context *sigma-finite-subalgebra*
begin

lemma *borel-measurable-cond-exp'[measurable]: cond-exp M F f ∈ borel-measurable M*

by (metis cond-exp-def someI has-cond-exp-def borel-measurable-const subalg measurable-from-subalg)

lemma *cond-exp-null:*

assumes $\nexists g.\ has_cond_exp\ M\ F\ f\ g$

shows $cond_exp\ M\ F\ f = (\lambda\cdot.\ 0)$

unfolding *cond-exp-def* **using** *assms* **by** *argo*

We state the tower property of the conditional expectation in terms of the predicate *has-cond-exp*.

lemma *has-cond-exp-nested-subalg:*

fixes $f :: 'a \Rightarrow 'b::\{second-countable-topology, banach\}$

assumes *subalgebra G F has-cond-exp M F f h has-cond-exp M G f h'*

shows *has-cond-exp M F h' h*

by (intro has-cond-expI') (metis assms has-cond-expD in-mono subalgebra-def)+

The following lemma shows that the conditional expectation is unique as an element of L1, given that it exists.

lemma *has-cond-exp-charact:*

fixes $f :: 'a \Rightarrow 'b::\{second-countable-topology, banach\}$

assumes *has-cond-exp M F f g*

shows *has-cond-exp M F f (cond-exp M F f)*

$AE\ x\ in\ M.\ cond_exp\ M\ F\ f\ x = g\ x$

proof –

show *cond-exp: has-cond-exp M F f (cond-exp M F f)* **using** *assms someI cond-exp-def* **by** *metis*

let $?MF = restr_to_subalg\ M\ F$

interpret *sigma-finite-measure ?MF* **by** (rule *sigma-fin-subalg*)

{

fix A **assume** $A \in sets\ ?MF$

then have $[measurable]: A \in sets\ F$ **using** *sets-restr-to-subalg[OF subalg]* **by** *simp*

have $(\int x \in A.\ g\ x\ \partial ?MF) = (\int x \in A.\ g\ x\ \partial M)$ **using** *assms subalg* **by** (auto *simp add: integral-subalgebra2 set-lebesgue-integral-def dest!: has-cond-expD*)

also have $\dots = (\int x \in A.\ cond_exp\ M\ F\ f\ x\ \partial M)$ **using** *assms cond-exp* **by** (*simp add: has-cond-exp-def*)

also have $\dots = (\int x \in A.\ cond_exp\ M\ F\ f\ x\ \partial ?MF)$ **using** *subalg* **by** (auto *simp add: integral-subalgebra2 set-lebesgue-integral-def*)

finally have $(\int x \in A.\ g\ x\ \partial ?MF) = (\int x \in A.\ cond_exp\ M\ F\ f\ x\ \partial ?MF)$ **by** *simp*

```

}
hence  $AE\ x\ in\ ?MF.\ cond\text{-}exp\ M\ F\ f\ x = g\ x$  using cond-exp assms subalg by
(intro density-unique-banach, auto dest: has-cond-expD intro!: integrable-in-subalg)
then show  $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ f\ x = g\ x$  using AE-restr-to-subalg[OF
subalg] by simp
qed

```

corollary *cond-exp-charact*:

```

fixes  $f :: 'a \Rightarrow 'b :: \{second\text{-countable-topology}, banach\}$ 
assumes  $\bigwedge A. A \in sets\ F \implies (\int x \in A. f\ x\ \partial M) = (\int x \in A. g\ x\ \partial M)$ 
 $integrable\ M\ f$ 
 $integrable\ M\ g$ 
 $g \in borel\text{-measurable}\ F$ 
shows  $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ f\ x = g\ x$ 
by (intro has-cond-exp-charact has-cond-expI' assms) auto

```

Identity on F-measurable functions:

If an integrable function f is already F -measurable, then $cond\text{-}exp\ M\ F\ f = f$ μ -a.e. This is a corollary of the lemma on the characterization of *cond-exp*.

corollary *cond-exp-F-meas[intro, simp]*:

```

fixes  $f :: 'a \Rightarrow 'b :: \{second\text{-countable-topology}, banach\}$ 
assumes  $integrable\ M\ f$ 
 $f \in borel\text{-measurable}\ F$ 
shows  $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ f\ x = f\ x$ 
by (rule cond-exp-charact, auto intro: assms)

```

Congruence

lemma *has-cond-exp-cong*:

```

assumes  $integrable\ M\ f \bigwedge x. x \in space\ M \implies f\ x = g\ x$  has-cond-exp M F g h
shows has-cond-exp M F f h
proof (intro has-cond-expI'[OF - assms(1)])
fix  $A$  assume asm: A ∈ sets F
hence  $set\text{-lebesgue-integral}\ M\ A\ f = set\text{-lebesgue-integral}\ M\ A\ g$  by (intro set-lebesgue-integral-cong)
(meson assms(2) subalg in-mono subalgebra-def sets.sets-into-space subalgebra-def
subsetD) +
thus  $set\text{-lebesgue-integral}\ M\ A\ f = set\text{-lebesgue-integral}\ M\ A\ h$  using asm assms(3)
by (simp add: has-cond-exp-def)
qed (auto simp add: has-cond-expD[OF assms(3)])

```

lemma *cond-exp-cong*:

```

fixes  $f :: 'a \Rightarrow 'b :: \{second\text{-countable-topology}, banach\}$ 
assumes  $integrable\ M\ f\ integrable\ M\ g \bigwedge x. x \in space\ M \implies f\ x = g\ x$ 
shows  $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ f\ x = cond\text{-}exp\ M\ F\ g\ x$ 
proof (cases  $\exists h. has\text{-}cond\text{-}exp\ M\ F\ f\ h$ )
case True
then obtain  $h$  where  $h: has\text{-}cond\text{-}exp\ M\ F\ f\ h\ has\text{-}cond\text{-}exp\ M\ F\ g\ h$  using
has-cond-exp-cong assms by metis
show ?thesis using  $h[THEN\ has\text{-}cond\text{-}exp\text{-}character(2)]$  by fastforce

```

```

next
  case False
  moreover have  $\nexists h. \text{has-cond-exp } M F g h$  using False has-cond-exp-cong assms
by auto
  ultimately show ?thesis unfolding cond-exp-def by auto
qed

lemma has-cond-exp-cong-AE:
  assumes integrable  $M f$  AE  $x$  in  $M. f x = g x$  has-cond-exp  $M F g h$ 
  shows has-cond-exp  $M F f h$ 
  using assms(1,2) subalg subalgebra-def subset-iff
  by (intro has-cond-expI', subst set-lebesgue-integral-cong-AE[OF - assms(1)] THEN
    borel-measurable-integrable] borel-measurable-integrable(1)[OF has-cond-expD(2)[OF
    assms(3)]]])
    (fast intro: has-cond-expD[OF assms(3)] integrable-cong-AE-imp[OF - - AE-symmetric])+

lemma has-cond-exp-cong-AE':
  assumes  $h \in \text{borel-measurable } F$  AE  $x$  in  $M. h x = h' x$  has-cond-exp  $M F f h'$ 
  shows has-cond-exp  $M F f h$ 
  using assms(1, 2) subalg subalgebra-def subset-iff
  using AE-restr-to-subalg2[OF subalg assms(2)] measurable-from-subalg
  by (intro has-cond-expI', subst set-lebesgue-integral-cong-AE[OF - measurable-from-subalg(1,1)] OF
    subalg], OF - assms(1) has-cond-expD(4)[OF assms(3)]]])
    (fast intro: has-cond-expD[OF assms(3)] integrable-cong-AE-imp[OF - - AE-symmetric])+

lemma cond-exp-cong-AE:
  fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
  assumes integrable  $M f$  integrable  $M g$  AE  $x$  in  $M. f x = g x$ 
  shows AE  $x$  in  $M. \text{cond-exp } M F f x = \text{cond-exp } M F g x$ 
proof (cases  $\exists h. \text{has-cond-exp } M F f h$ )
  case True
  then obtain  $h$  where  $h: \text{has-cond-exp } M F f h$  has-cond-exp  $M F g h$  using
    has-cond-exp-cong-AE assms by (metis (mono-tags, lifting) eventually-mono)
  show ?thesis using  $h$  [THEN has-cond-exp-charact(2)] by fastforce
next
  case False
  moreover have  $\nexists h. \text{has-cond-exp } M F g h$  using False has-cond-exp-cong-AE
    assms by auto
  ultimately show ?thesis unfolding cond-exp-def by auto
qed

```

The conditional expectation operator on the reals, *real-cond-exp*, satisfies the conditions of the conditional expectation as we have defined it.

```

lemma has-cond-exp-real:
  fixes  $f :: 'a \Rightarrow \text{real}$ 
  assumes integrable  $M f$ 
  shows has-cond-exp  $M F f$  (real-cond-exp  $M F f$ )
  by (intro has-cond-expI', auto intro!: real-cond-exp-intA assms)

```

```

lemma cond-exp-real[intro]:
  fixes  $f :: 'a \Rightarrow \text{real}$ 
  assumes integrable  $M f$ 
  shows  $AE\ x\ in\ M.\ \text{cond-exp}\ M\ F\ f\ x = \text{real-cond-exp}\ M\ F\ f\ x$ 
  using has-cond-exp-charact has-cond-exp-real assms by blast

lemma cond-exp-cmult:
  fixes  $f :: 'a \Rightarrow \text{real}$ 
  assumes integrable  $M f$ 
  shows  $AE\ x\ in\ M.\ \text{cond-exp}\ M\ F\ (\lambda x.\ c * f\ x)\ x = c * \text{cond-exp}\ M\ F\ f\ x$ 
  using real-cond-exp-cmult[OF assms(1), of  $c$ ] assms(1)[THEN cond-exp-real]
assms(1)[THEN integrable-mult-right, THEN cond-exp-real, of  $c$ ] by fastforce

```

4.1 Existence

Showing the existence is a bit involved. Specifically, what we aim to show is that $\text{has-cond-exp}\ M\ F\ f\ (\text{cond-exp}\ M\ F\ f)$ holds for any Bochner-integrable f . We will employ the standard machinery of measure theory. First, we will prove existence for indicator functions. Then we will extend our proof by linearity to simple functions. Finally we use a limiting argument to show that the conditional expectation exists for all Bochner-integrable functions.

Indicator functions

```

lemma has-cond-exp-indicator:
  assumes  $A \in \text{sets}\ M\ \text{emeasure}\ M\ A < \infty$ 
  shows  $\text{has-cond-exp}\ M\ F\ (\lambda x.\ \text{indicat-real}\ A\ x *_{\mathbb{R}} y)\ (\lambda x.\ \text{real-cond-exp}\ M\ F\ (\text{indicator}\ A)\ x *_{\mathbb{R}} y)$ 
proof (intro has-cond-expI', goal-cases)
  case 1  $B$ 
  have  $\int_{x \in B} (\text{indicat-real}\ A\ x *_{\mathbb{R}} y)\ \partial M = (\int_{x \in B} \text{indicat-real}\ A\ x\ \partial M) *_{\mathbb{R}} y$ 
  using assms by (intro set-integral-scaleR-left, meson 1 in-mono subalg subalgebra-def, blast)
  also have  $\dots = (\int_{x \in B} \text{real-cond-exp}\ M\ F\ (\text{indicator}\ A)\ x\ \partial M) *_{\mathbb{R}} y$ 
  using 1 assms by (subst real-cond-exp-intA, auto)
  also have  $\dots = \int_{x \in B} (\text{real-cond-exp}\ M\ F\ (\text{indicator}\ A)\ x *_{\mathbb{R}} y)\ \partial M$ 
  using assms by (intro set-integral-scaleR-left[symmetric], meson 1 in-mono subalg subalgebra-def, blast)
  finally show ?case .
next
  case 2
  show ?case using integrable-scaleR-left integrable-real-indicator assms by blast
next
  case 3
  show ?case using assms by (intro integrable-scaleR-left, intro real-cond-exp-int, blast+)
next
  case 4
  show ?case by (intro borel-measurable-scaleR, intro Conditional-Expectation.borel-measurable-cond-exp,

```


simp)
qed

lemma *cond-exp-indicator*[intro]:
fixes $y :: 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes [measurable]: $A \in \text{sets } M \text{ emeasure } M A < \infty$
shows $AE\ x \text{ in } M. \text{cond-exp } M\ F\ (\lambda x. \text{indicat-real } A\ x *_R y)\ x = \text{cond-exp } M\ F$
 $(\text{indicator } A)\ x *_R y$
proof –
have $AE\ x \text{ in } M. \text{cond-exp } M\ F\ (\lambda x. \text{indicat-real } A\ x *_R y)\ x = \text{real-cond-exp } M\ F$
 $(\text{indicator } A)\ x *_R y$ **using** *has-cond-exp-indicator*[OF assms] *has-cond-exp-charact*
by *blast*
thus ?thesis **using** *cond-exp-real*[OF *integrable-real-indicator*, OF assms] **by** *fast-force*
qed

Addition

lemma *has-cond-exp-add*:
fixes $f\ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *has-cond-exp* $M\ F\ f\ f'$ *has-cond-exp* $M\ F\ g\ g'$
shows *has-cond-exp* $M\ F\ (\lambda x. f\ x + g\ x)\ (\lambda x. f'\ x + g'\ x)$
proof (intro *has-cond-expI'*, goal-cases)
case (1 A)
have $\int_{x \in A}. (f\ x + g\ x) \partial M = (\int_{x \in A}. f\ x \partial M) + (\int_{x \in A}. g\ x \partial M)$ **using**
assms[THEN *has-cond-expD*(2)] *subalg* 1 **by** (intro *set-integral-add*(2), auto *simp*
add: subalgebra-def set-integrable-def intro: integrable-mult-indicator)
also have $\dots = (\int_{x \in A}. f'\ x \partial M) + (\int_{x \in A}. g'\ x \partial M)$ **using** *assms*[THEN
has-cond-expD(1)[OF - 1]] **by** *argo*
also have $\dots = \int_{x \in A}. (f'\ x + g'\ x) \partial M$ **using** *assms*[THEN *has-cond-expD*(3)]
subalg 1 **by** (intro *set-integral-add*(2)[*symmetric*], auto *simp add: subalgebra-def*
set-integrable-def intro: integrable-mult-indicator)
finally show ?case .
next
case 2
show ?case **by** (metis *Bochner-Integration.integrable-add assms has-cond-expD*(2))
next
case 3
show ?case **by** (metis *Bochner-Integration.integrable-add assms has-cond-expD*(3))
next
case 4
show ?case **using** *assms borel-measurable-add has-cond-expD*(4) **by** *blast*
qed

lemma *has-cond-exp-scaleR-right*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *has-cond-exp* $M\ F\ f\ f'$
shows *has-cond-exp* $M\ F\ (\lambda x. c *_R f\ x)\ (\lambda x. c *_R f'\ x)$
using *has-cond-expD*[OF assms] **by** (intro *has-cond-expI'*, auto)

```

lemma cond-exp-scaleR-right:
  fixes  $f :: 'a \Rightarrow 'b::\{\text{second-countable-topology}, \text{banach}\}$ 
  assumes integrable  $M f$ 
  shows  $\text{AE } x \text{ in } M. \text{cond-exp } M F (\lambda x. c *_R f x) x = c *_R \text{cond-exp } M F f x$ 
proof (cases  $\exists f'. \text{has-cond-exp } M F f f'$ )
  case True
    then show ?thesis using assms has-cond-exp-charact has-cond-exp-scaleR-right
  by metis
next
  case False
  show ?thesis
  proof (cases  $c = 0$ )
  case True
    then show ?thesis by simp
  next
  case c-nonzero: False
  have  $\nexists f'. \text{has-cond-exp } M F (\lambda x. c *_R f x) f'$ 
  proof (standard, goal-cases)
  case 1
    then obtain  $f'$  where  $f': \text{has-cond-exp } M F (\lambda x. c *_R f x) f'$  by blast
    have  $\text{has-cond-exp } M F f (\lambda x. \text{inverse } c *_R f' x)$  using has-cond-expD[OF
 $f'$ ] divideR-right[OF c-nonzero] assms by (intro has-cond-expI', auto)
    then show ?case using False by blast
  qed
  then show ?thesis using cond-exp-null[OF False] cond-exp-null by force
  qed
qed

```

```

lemma cond-exp-uminus:
  fixes  $f :: 'a \Rightarrow 'b::\{\text{second-countable-topology}, \text{banach}\}$ 
  assumes integrable  $M f$ 
  shows  $\text{AE } x \text{ in } M. \text{cond-exp } M F (\lambda x. - f x) x = - \text{cond-exp } M F f x$ 
  using cond-exp-scaleR-right[OF assms, of -1] by force

```

Together with the induction scheme *integrable-simple-function-induct*, we can show that the conditional expectation of an integrable simple function exists.

```

corollary has-cond-exp-simple:
  fixes  $f :: 'a \Rightarrow 'b::\{\text{second-countable-topology}, \text{banach}\}$ 
  assumes simple-function  $M f$  emeasure  $M \{y \in \text{space } M. f y \neq 0\} \neq \infty$ 
  shows  $\text{has-cond-exp } M F f (\text{cond-exp } M F f)$ 
  using assms
proof (induction rule: integrable-simple-function-induct)
  case (cong  $f g$ )
    then show ?case using has-cond-exp-cong by (metis (no-types, opaque-lifting)
Bochner-Integration.integrable-cong has-cond-expD(2) has-cond-exp-charact(1))
  next
  case (indicator  $A y$ )
  then show ?case using has-cond-exp-charact[OF has-cond-exp-indicator] by fast

```

```

next
  case (add u v)
  then show ?case using has-cond-exp-add has-cond-exp-charact(1) by blast
qed

```

Now comes the most difficult part. Given a convergent sequence of integrable simple functions s , we must show that the sequence $\lambda n. \text{cond-exp } M F (s\ n)$ is also convergent. Furthermore, we must show that this limit satisfies the properties of a conditional expectation. Unfortunately, we will only be able to show that this sequence converges in the L1-norm. Luckily, this is enough to show that the operator $\text{cond-exp } M F$ preserves limits as a function from L1 to L1.

In anticipation of this result, we show that the conditional expectation operator is a contraction for simple functions. We first reformulate the lemma *real-cond-exp-abs*, which shows the statement for real-valued functions, using our definitions. Then we show the statement for simple functions via induction.

lemma *cond-exp-contraction-real*:

```

  fixes f :: 'a  $\Rightarrow$  real
  assumes integrable[measurable]: integrable M f
  shows AE x in M. norm (cond-exp M F f x)  $\leq$  cond-exp M F ( $\lambda x. \text{norm } (f\ x)$ ) x
proof-
  have int: integrable M ( $\lambda x. \text{norm } (f\ x)$ ) using assms by blast
  have *: AE x in M.  $0 \leq \text{cond-exp } M F (\lambda x. \text{norm } (f\ x))\ x$  using cond-exp-real[THEN
    AE-symmetric, OF integrable-norm[OF integrable]] real-cond-exp-ge-c[OF integrable-norm[OF
    integrable], of 0] norm-ge-zero by fastforce
  have **:  $A \in \text{sets } F \implies \int x \in A. |f\ x| \partial M = \int x \in A. \text{real-cond-exp } M F (\lambda x. \text{norm } (f\ x))\ x \partial M$  for A
    unfolding real-norm-def using assms integrable-abs real-cond-exp-intA by blast

  have norm-int:  $A \in \text{sets } F \implies (\int x \in A. |f\ x| \partial M) = (\int^+ x \in A. |f\ x| \partial M)$  for A
    using assms by (intro nn-set-integral-eq-set-integral[symmetric], blast, fastforce) (meson subalg subalgebra-def subsetD)

  have AE x in M. real-cond-exp M F ( $\lambda x. \text{norm } (f\ x)$ ) x  $\geq 0$  using int real-cond-exp-ge-c
    by force
  hence cond-exp-norm-int:  $A \in \text{sets } F \implies (\int x \in A. \text{real-cond-exp } M F (\lambda x. \text{norm } (f\ x))\ x \partial M) = (\int^+ x \in A. \text{real-cond-exp } M F (\lambda x. \text{norm } (f\ x))\ x \partial M)$  for A
    using assms by (intro nn-set-integral-eq-set-integral[symmetric], blast, fastforce) (meson subalg subalgebra-def subsetD)

  have  $A \in \text{sets } F \implies \int^+ x \in A. |f\ x| \partial M = \int^+ x \in A. \text{real-cond-exp } M F (\lambda x. \text{norm } (f\ x))\ x \partial M$  for A
    using ** norm-int cond-exp-norm-int by (auto simp add: nn-integral-set-ennreal)
  moreover have ( $\lambda x. \text{ennreal } |f\ x|$ )  $\in \text{borel-measurable } M$  by measurable
  moreover have ( $\lambda x. \text{ennreal } (\text{real-cond-exp } M F (\lambda x. \text{norm } (f\ x))\ x)$ )  $\in \text{borel-measurable } F$ 
    by measurable

```

ultimately have $AE\ x\ in\ M. nn\text{-}cond\text{-}exp\ M\ F\ (\lambda x. ennreal\ |f\ x|)\ x = real\text{-}cond\text{-}exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x$ **by** (intro nn-cond-exp-charact[THEN AE-symmetric], auto)
hence $AE\ x\ in\ M. nn\text{-}cond\text{-}exp\ M\ F\ (\lambda x. ennreal\ |f\ x|)\ x \leq cond\text{-}exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x$ **using** cond-exp-real[OF int] **by** force
moreover have $AE\ x\ in\ M. |real\text{-}cond\text{-}exp\ M\ F\ f\ x| = norm\ (cond\text{-}exp\ M\ F\ f\ x)$
unfolding real-norm-def **using** cond-exp-real[OF assms] * **by** force
ultimately have $AE\ x\ in\ M. ennreal\ (norm\ (cond\text{-}exp\ M\ F\ f\ x)) \leq cond\text{-}exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x$ **using** real-cond-exp-abs[OF assms[THEN borel-measurable-integrable]] **by** fastforce
hence $AE\ x\ in\ M. enn2real\ (ennreal\ (norm\ (cond\text{-}exp\ M\ F\ f\ x))) \leq enn2real\ (cond\text{-}exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x)$ **using** ennreal-le-iff2 **by** force
thus ?thesis **using** * **by** fastforce
qed

lemma cond-exp-contraction-simple:

fixes $f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach\}$
assumes simple-function $M\ f\ emeasure\ M\ \{y \in space\ M. f\ y \neq 0\} \neq \infty$
shows $AE\ x\ in\ M. norm\ (cond\text{-}exp\ M\ F\ f\ x) \leq cond\text{-}exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x$ **using** assms
proof (induction rule: integrable-simple-function-induct)
case (cong f g)
hence $ae: AE\ x\ in\ M. f\ x = g\ x$ **by** blast
hence $AE\ x\ in\ M. cond\text{-}exp\ M\ F\ f\ x = cond\text{-}exp\ M\ F\ g\ x$ **using** cong has-cond-exp-simple **by** (subst cond-exp-cong-AE) (auto intro!: has-cond-expD(2))
hence $AE\ x\ in\ M. norm\ (cond\text{-}exp\ M\ F\ f\ x) = norm\ (cond\text{-}exp\ M\ F\ g\ x)$ **by** force
moreover have $AE\ x\ in\ M. cond\text{-}exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x = cond\text{-}exp\ M\ F\ (\lambda x. norm\ (g\ x))\ x$ **using** ae cong has-cond-exp-simple **by** (subst cond-exp-cong-AE) (auto dest: has-cond-expD)
ultimately show ?case **using** cong(6) **by** fastforce
next
case (indicator A y)
hence $AE\ x\ in\ M. cond\text{-}exp\ M\ F\ (\lambda a. indicator\ A\ a *_{\mathbb{R}} y)\ x = cond\text{-}exp\ M\ F\ (indicator\ A)\ x *_{\mathbb{R}} y$ **by** blast
hence *: $AE\ x\ in\ M. norm\ (cond\text{-}exp\ M\ F\ (\lambda a. indicat\text{-}real\ A\ a *_{\mathbb{R}} y)\ x) \leq norm\ y * cond\text{-}exp\ M\ F\ (\lambda x. norm\ (indicat\text{-}real\ A\ x))\ x$ **using** cond-exp-contraction-real[OF integrable-real-indicator, OF indicator] **by** fastforce

have $AE\ x\ in\ M. norm\ y * cond\text{-}exp\ M\ F\ (\lambda x. norm\ (indicat\text{-}real\ A\ x))\ x = norm\ y * real\text{-}cond\text{-}exp\ M\ F\ (\lambda x. norm\ (indicat\text{-}real\ A\ x))\ x$ **using** cond-exp-real[OF integrable-real-indicator, OF indicator] **by** fastforce

moreover have $AE\ x\ in\ M. cond\text{-}exp\ M\ F\ (\lambda x. norm\ y * norm\ (indicat\text{-}real\ A\ x))\ x = real\text{-}cond\text{-}exp\ M\ F\ (\lambda x. norm\ y * norm\ (indicat\text{-}real\ A\ x))\ x$ **using** indicator **by** (intro cond-exp-real, auto)

ultimately have $AE\ x\ in\ M. norm\ y * cond\text{-}exp\ M\ F\ (\lambda x. norm\ (indicat\text{-}real\ A\ x))\ x = cond\text{-}exp\ M\ F\ (\lambda x. norm\ y * norm\ (indicat\text{-}real\ A\ x))\ x$ **using** real-cond-exp-cmult[of $\lambda x. norm\ (indicat\text{-}real\ A\ x)\ norm\ y$] indicator **by** fastforce

moreover have $(\lambda x. norm\ y * norm\ (indicat\text{-}real\ A\ x)) = (\lambda x. norm\ (indicat\text{-}real\ A\ x) * norm\ y)$

$A \ x *_R \ y))$ **by force**
ultimately show $?case$ **using** $*$ **by force**
next
case $(add \ u \ v)$
have $AE \ x \ in \ M. \ norm \ (cond-exp \ M \ F \ (\lambda a. \ u \ a + v \ a) \ x) = norm \ (cond-exp \ M \ F \ u \ x + cond-exp \ M \ F \ v \ x)$ **using** $has-cond-exp-charact(2)[OF \ has-cond-exp-add, OF \ has-cond-exp-simple(1,1), OF \ add(1,2,3,4)]$ **by fastforce**
moreover have $AE \ x \ in \ M. \ norm \ (cond-exp \ M \ F \ u \ x + cond-exp \ M \ F \ v \ x) \leq norm \ (cond-exp \ M \ F \ u \ x) + norm \ (cond-exp \ M \ F \ v \ x)$ **using** $norm-triangle-ineq$ **by blast**
moreover have $AE \ x \ in \ M. \ norm \ (cond-exp \ M \ F \ u \ x) + norm \ (cond-exp \ M \ F \ v \ x) \leq cond-exp \ M \ F \ (\lambda x. \ norm \ (u \ x)) \ x + cond-exp \ M \ F \ (\lambda x. \ norm \ (v \ x)) \ x$ **using** $add(6,7)$ **by fastforce**
moreover have $AE \ x \ in \ M. \ cond-exp \ M \ F \ (\lambda x. \ norm \ (u \ x)) \ x + cond-exp \ M \ F \ (\lambda x. \ norm \ (v \ x)) \ x = cond-exp \ M \ F \ (\lambda x. \ norm \ (u \ x) + norm \ (v \ x)) \ x$ **using** $integrable-simple-function[OF \ add(1,2)] \ integrable-simple-function[OF \ add(3,4)]$ **by** $(intro \ has-cond-exp-charact(2)[OF \ has-cond-exp-add[OF \ has-cond-exp-charact(1,1)], THEN \ AE-symmetric], auto \ intro: \ has-cond-exp-real)$
moreover have $AE \ x \ in \ M. \ cond-exp \ M \ F \ (\lambda x. \ norm \ (u \ x) + norm \ (v \ x)) \ x = cond-exp \ M \ F \ (\lambda x. \ norm \ (u \ x + v \ x)) \ x$ **using** $add(5) \ integrable-simple-function[OF \ add(1,2)] \ integrable-simple-function[OF \ add(3,4)]$ **by** $(intro \ cond-exp-cong, auto)$
ultimately show $?case$ **by force**
qed

lemma $has-cond-exp-simple-lim$:

fixes $f :: 'a \Rightarrow 'b :: \{second-countable-topology, banach\}$
assumes $integrable[measurable]: \ integrable \ M \ f$
and $\bigwedge i. \ simple-function \ M \ (s \ i)$
and $\bigwedge i. \ emeasure \ M \ \{y \in space \ M. \ s \ i \ y \neq 0\} \neq \infty$
and $\bigwedge x. \ x \in space \ M \implies (\lambda i. \ s \ i \ x) \longrightarrow f \ x$
and $\bigwedge x \ i. \ x \in space \ M \implies norm \ (s \ i \ x) \leq 2 * norm \ (f \ x)$
obtains r
where $strict-mono \ r \ has-cond-exp \ M \ F \ f \ (\lambda x. \ lim \ (\lambda i. \ cond-exp \ M \ F \ (s \ (r \ i)) \ x))$
 $AE \ x \ in \ M. \ convergent \ (\lambda i. \ cond-exp \ M \ F \ (s \ (r \ i)) \ x)$

proof –

have $[measurable]: (s \ i) \in borel-measurable \ M$ **for** i **using** $assms(2)$ **by** $(simp \ add: \ borel-measurable-simple-function)$
have $integrable-s: \ integrable \ M \ (\lambda x. \ s \ i \ x)$ **for** i **using** $assms \ integrable-simple-function$ **by blast**
have $integrable-4f: \ integrable \ M \ (\lambda x. \ 4 * norm \ (f \ x))$ **using** $assms(1)$ **by simp**
have $integrable-2f: \ integrable \ M \ (\lambda x. \ 2 * norm \ (f \ x))$ **using** $assms(1)$ **by simp**
have $integrable-2-cond-exp-norm-f: \ integrable \ M \ (\lambda x. \ 2 * cond-exp \ M \ F \ (\lambda x. \ norm \ (f \ x)) \ x)$ **by fast**

have $emeasure \ M \ \{y \in space \ M. \ s \ i \ y - s \ j \ y \neq 0\} \leq emeasure \ M \ \{y \in space \ M. \ s \ i \ y \neq 0\} + emeasure \ M \ \{y \in space \ M. \ s \ j \ y \neq 0\}$ **for** $i \ j$ **using** $simple-functionD(2)[OF \ assms(2)]$ **by** $(intro \ order-trans[OF \ emeasure-mono \ emeasure-subadditive], auto)$

hence *fin-sup*: $\text{emeasure } M \{y \in \text{space } M. s \ i \ y - s \ j \ y \neq 0\} \neq \infty$ **for** $i \ j$ **using** *assms(3)* **by** (*metis (mono-tags) ennreal-add-eq-top linorder-not-less top.not-eq-extremum infinity-ennreal-def*)

have $\text{emeasure } M \{y \in \text{space } M. \text{norm } (s \ i \ y - s \ j \ y) \neq 0\} \leq \text{emeasure } M \{y \in \text{space } M. s \ i \ y \neq 0\} + \text{emeasure } M \{y \in \text{space } M. s \ j \ y \neq 0\}$ **for** $i \ j$ **using** *simple-functionD(2)[OF assms(2)]* **by** (*intro order-trans[OF emeasure-mono emeasure-subadditive], auto*)

hence *fin-sup-norm*: $\text{emeasure } M \{y \in \text{space } M. \text{norm } (s \ i \ y - s \ j \ y) \neq 0\} \neq \infty$ **for** $i \ j$ **using** *assms(3)* **by** (*metis (mono-tags) ennreal-add-eq-top linorder-not-less top.not-eq-extremum infinity-ennreal-def*)

have *Cauchy*: $\text{Cauchy } (\lambda n. s \ n \ x)$ **if** $x \in \text{space } M$ **for** x **using** *assms(4) LIM-SEQ-imp-Cauchy* **that** **by** *blast*

hence *bounded-range-s*: $\text{bounded } (\text{range } (\lambda n. s \ n \ x))$ **if** $x \in \text{space } M$ **for** x **using** *that cauchy-imp-bounded* **by** *fast*

Since the sequence $\lambda n. s \ n \ x$ is Cauchy for almost all x , we know that the diameter tends to zero almost everywhere.

Dominated convergence tells us that the integral of the diameter also converges to zero.

have $\text{AE } x \text{ in } M. (\lambda n. \text{diameter } \{s \ i \ x \mid i. n \leq i\}) \longrightarrow 0$ **using** *Cauchy cauchy-iff-diameter-tends-to-zero-and-bounded* **by** *fast*

moreover **have** $(\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\}) \in \text{borel-measurable } M$ **for** n **using** *bounded-range-s borel-measurable-diameter* **by** *measurable*

moreover **have** $\text{AE } x \text{ in } M. \text{norm } (\text{diameter } \{s \ i \ x \mid i. n \leq i\}) \leq 4 * \text{norm } (f \ x)$ **for** n

proof –

{

fix x **assume** $x: x \in \text{space } M$

have $\text{diameter } \{s \ i \ x \mid i. n \leq i\} \leq 2 * \text{norm } (f \ x) + 2 * \text{norm } (f \ x)$ **by** (*intro diameter-le, blast, subst dist-norm[symmetric], intro dist-triangle3[THEN order-trans, of 0], intro add-mono*) (*auto intro: assms(5)[OF x]*)

hence $\text{norm } (\text{diameter } \{s \ i \ x \mid i. n \leq i\}) \leq 4 * \text{norm } (f \ x)$ **using** *diameter-ge-0[OF bounded-subset[OF bounded-range-s], OF x, of {s i x | i. n ≤ i}]* **by** *force*

}

thus *?thesis* **by** *fast*

qed

ultimately **have** *diameter-tendsto-zero*: $(\lambda n. \text{LINT } x | M. \text{diameter } \{s \ i \ x \mid i. n \leq i\}) \longrightarrow 0$ **by** (*intro integral-dominated-convergence[OF borel-measurable-const[of 0] - integrable-4f, simplified]*) (*fast+*)

have *diameter-integrable*: $\text{integrable } M (\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\})$ **for** n **using** *assms(1,5)*

by (*intro integrable-bound-diameter[OF bounded-range-s integrable-2f], auto*)

have *dist-integrable*: $\text{integrable } M (\lambda x. \text{dist } (s \ i \ x) (s \ j \ x))$ **for** $i \ j$ **using** *assms(5)*

dist-triangle3[*of s i - 0, THEN order-trans, OF add-mono, of - 2 * norm (f -)*]
by (*intro Bochner-Integration.integrable-bound*[*OF integrable-4f*]) *fastforce*+

Since *cond-exp M F* is a contraction for simple functions, the following sequence of integral values is also Cauchy.

This follows, since the distance between the terms of this sequence are always less than or equal to the diameter, which itself converges to zero.

Hence, we obtain a subsequence which is Cauchy almost everywhere.

have $\exists N. \forall i \geq N. \forall j \geq N. \text{LINT } x | M. \text{norm } (\text{cond-exp } M F (s i) x - \text{cond-exp } M F (s j) x) < e$ **if** *e-pos*: *e > 0* **for** *e*

proof -

obtain *N* **where** *: *LINT x | M. diameter* $\{s i x \mid i. n \leq i\} < e$ **if** *n ≥ N* **for** *n* **using** *that order-tendsto-iff*[*THEN iffD1, OF diameter-tendsto-zero, unfolded eventually-sequentially*] *e-pos* **by** *presburger*

{

fix *i j x* **assume** *asm*: $i \geq N j \geq N x \in \text{space } M$

have *case-prod dist* ' ($\{s i x \mid i. N \leq i\} \times \{s i x \mid i. N \leq i\}$) = *case-prod* ($\lambda i j. \text{dist } (s i x) (s j x)$) ' ($\{N..\} \times \{N..\}$) **by** *fast*

hence *diameter* $\{s i x \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s i x) (s j x))$ **unfolding** *diameter-def* **by** *auto*

moreover **have** $(\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s i x) (s j x)) \geq \text{dist } (s i x) (s j x)$ **using** *asm bounded-imp-bdd-above*[*OF bounded-imp-dist-bounded, OF bounded-range-s*] **by** (*intro cSup-upper, auto*)

ultimately **have** *diameter* $\{s i x \mid i. N \leq i\} \geq \text{dist } (s i x) (s j x)$ **by** *presburger*

}

hence *LINT x | M. dist* $(s i x) (s j x) < e$ **if** $i \geq N j \geq N$ **for** *i j* **using** *that ** **by** (*intro integral-mono*[*OF dist-integrable diameter-integrable, THEN order.strict-trans1*], *blast*+))

moreover **have** *LINT x | M. norm* $(\text{cond-exp } M F (s i) x - \text{cond-exp } M F (s j) x) \leq \text{LINT } x | M. \text{dist } (s i x) (s j x)$ **for** *i j*

proof -

have *LINT x | M. norm* $(\text{cond-exp } M F (s i) x - \text{cond-exp } M F (s j) x) = \text{LINT } x | M. \text{norm } (\text{cond-exp } M F (s i) x + - 1 *_R \text{cond-exp } M F (s j) x)$ **unfolding** *dist-norm* **by** *simp*

also **have** ... = *LINT x | M. norm* $(\text{cond-exp } M F (\lambda x. s i x - s j x) x)$ **using** *has-cond-exp-charact(2)*[*OF has-cond-exp-add*[*OF - has-cond-exp-scaleR-right, OF has-cond-exp-charact(1,1), OF has-cond-exp-simple(1,1)*][*OF assms(2,3)*]], *THEN AE-symmetric, of i - 1 j*] **by** (*intro integral-cong-AE*) *force*+

also **have** ... $\leq \text{LINT } x | M. \text{cond-exp } M F (\lambda x. \text{norm } (s i x - s j x)) x$ **using** *cond-exp-contraction-simple*[*OF - fin-sup, of i j*] *integrable-cond-exp assms(2)* **by** (*intro integral-mono-AE, fast*+))

also **have** ... = *LINT x | M. norm* $(s i x - s j x)$ **unfolding** *set-integral-space(1)*[*OF integrable-cond-exp, symmetric*] *set-integral-space*[*OF dist-integrable*[*unfolded dist-norm*], *symmetric*] **by** (*intro has-cond-expD(1)*[*OF has-cond-exp-simple*[*OF - fin-sup-norm*], *symmetric*]) (*metis assms(2) simple-function-compose1 simple-function-diff, metis sets.top subalg subalgebra-def*)

finally show *?thesis unfolding dist-norm* .
qed
ultimately show *?thesis using order.strict-trans1 by meson*
qed
then obtain *r where strict-mono-r: strict-mono r and AE-Cauchy: AE x in M. Cauchy (λi. cond-exp M F (s (r i)) x)*
by *(rule cauchy-L1-AE-cauchy-subseq[OF integrable-cond-exp], auto)*
hence *ae-lim-cond-exp: AE x in M. (λn. cond-exp M F (s (r n)) x) ⟶ lim (λn. cond-exp M F (s (r n)) x) using Cauchy-convergent-iff convergent-LIMSEQ-iff*
by *fastforce*

Now that we have a candidate for the conditional expectation, we must show that it actually has the required properties.

Dominated convergence shows that this limit is indeed integrable.

Here, we again use the fact that conditional expectation is a contraction on simple functions.

have *cond-exp-bounded: AE x in M. norm (cond-exp M F (s (r n)) x) ≤ cond-exp M F (λx. 2 * norm (f x)) x for n*
proof –
have *AE x in M. norm (cond-exp M F (s (r n)) x) ≤ cond-exp M F (λx. norm (s (r n) x)) x* **by** *(rule cond-exp-contraction-simple[OF assms(2,3)])*
moreover have *AE x in M. real-cond-exp M F (λx. norm (s (r n) x)) x ≤ real-cond-exp M F (λx. 2 * norm (f x)) x* **using** *integrable-s integrable-2f assms(5)*
by *(intro real-cond-exp-mono, auto)*
ultimately show *?thesis using cond-exp-real[OF integrable-norm, OF integrable-s, of r n] cond-exp-real[OF integrable-2f] by force*
qed
have *lim-integrable: integrable M (λx. lim (λi. cond-exp M F (s (r i)) x))*
by *(intro integrable-dominated-convergence[OF - borel-measurable-cond-exp' integrable-cond-exp ae-lim-cond-exp cond-exp-bounded], simp)*

Moreover, we can use the DCT twice to show that the conditional expectation property holds, i.e. the value of the integral of the candidate, agrees with f on sets $A \in \text{sets } F$.

{
fix *A assume A-in-sets-F: A ∈ sets F*
have *AE x in M. norm (indicator A x *_R cond-exp M F (s (r n)) x) ≤ cond-exp M F (λx. 2 * norm (f x)) x* **for** *n*
proof –
have *AE x in M. norm (indicator A x *_R cond-exp M F (s (r n)) x) ≤ norm (cond-exp M F (s (r n)) x)* **unfolding** *indicator-def* **by** *simp*
thus *?thesis using cond-exp-bounded[of n] by force*
qed
hence *lim-cond-exp-int: (λn. LINT x:A|M. cond-exp M F (s (r n)) x) ⟶ LINT x:A|M. lim (λn. cond-exp M F (s (r n)) x)*
using *ae-lim-cond-exp measurable-from-subalg[OF subalg borel-measurable-indicator, OF A-in-sets-F] cond-exp-bounded*

unfolding *set-lebesgue-integral-def*
by (*intro integral-dominated-convergence*[*OF borel-measurable-scaleR borel-measurable-scaleR integrable-cond-exp*]) (*fastforce simp add: tendsto-scaleR*) +

have $AE\ x\ in\ M.\ norm\ (indicator\ A\ x\ *_R\ s\ (r\ n)\ x) \leq 2 * norm\ (f\ x)$ **for** n
proof –
have $AE\ x\ in\ M.\ norm\ (indicator\ A\ x\ *_R\ s\ (r\ n)\ x) \leq norm\ (s\ (r\ n)\ x)$
unfolding *indicator-def* **by** *simp*
thus *?thesis* **using** *assms(5)[of - r n]* **by** *fastforce*
qed
hence $lim\text{-}s\text{-}int: (\lambda n.\ LINT\ x:A|M.\ s\ (r\ n)\ x) \longrightarrow LINT\ x:A|M.\ f\ x$
using *measurable-from-subalg*[*OF subalg borel-measurable-indicator, OF A-in-sets-F*]
LIMSEQ-subseq-LIMSEQ[*OF assms(4) strict-mono-r*] *assms(5)*
unfolding *set-lebesgue-integral-def comp-def*
by (*intro integral-dominated-convergence*[*OF borel-measurable-scaleR borel-measurable-scaleR integrable-2f*]) (*fastforce simp add: tendsto-scaleR*) +

have $LINT\ x:A|M.\ lim\ (\lambda n.\ cond\text{-}exp\ M\ F\ (s\ (r\ n))\ x) = lim\ (\lambda n.\ LINT\ x:A|M.\ cond\text{-}exp\ M\ F\ (s\ (r\ n))\ x)$ **using** *limI*[*OF lim-cond-exp-int*] **by** *argo*
also have $\dots = lim\ (\lambda n.\ LINT\ x:A|M.\ s\ (r\ n)\ x)$ **using** *has-cond-expD(1)*[*OF has-cond-exp-simple*][*OF assms(2,3)*] *A-in-sets-F, symmetric*] **by** *presburger*
also have $\dots = LINT\ x:A|M.\ f\ x$ **using** *limI*[*OF lim-s-int*] **by** *argo*
finally have $LINT\ x:A|M.\ lim\ (\lambda n.\ cond\text{-}exp\ M\ F\ (s\ (r\ n))\ x) = LINT\ x:A|M.\ f\ x$.
}

Putting it all together, we have the statement we are looking for.

hence $has\text{-}cond\text{-}exp\ M\ F\ f\ (\lambda x.\ lim\ (\lambda i.\ cond\text{-}exp\ M\ F\ (s\ (r\ i))\ x))$ **using** *assms(1) lim-integrable* **by** (*intro has-cond-expI', auto*)
thus *thesis* **using** *AE-Cauchy Cauchy-convergent strict-mono-r* **by** (*auto intro!: that*)
qed

Now, we can show that the conditional expectation is well-defined for all integrable functions.

corollary *has-cond-expI*:

fixes $f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach\}$
assumes *integrable M f*
shows $has\text{-}cond\text{-}exp\ M\ F\ f\ (cond\text{-}exp\ M\ F\ f)$

proof –

obtain s **where** $s\text{-}is: \bigwedge i.\ simple\text{-}function\ M\ (s\ i) \wedge i.\ emeasure\ M\ \{y \in space\ M.\ s\ i\ y \neq 0\} \neq \infty \wedge x \in space\ M \implies (\lambda i.\ s\ i\ x) \longrightarrow f\ x \wedge x \in space\ M \implies norm\ (s\ i\ x) \leq 2 * norm\ (f\ x)$ **using** *integrable-implies-simple-function-sequence*[*OF assms*] **by** *blast*

show *?thesis* **using** *has-cond-exp-simple-lim*[*OF assms s-is*] *has-cond-exp-charact(1)*
by *metis*
qed

4.2 Properties

The defining property of the conditional expectation now always holds, given that the function f is integrable.

lemma *cond-exp-set-integral*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f A \in \text{sets } F$
shows $(\int x \in A. f x \partial M) = (\int x \in A. \text{cond-exp } M F f x \partial M)$
using *has-cond-expD(1)* [*OF has-cond-expI, OF assms*] **by** *argo*

The following property of the conditional expectation is called the "Tower Property".

lemma *cond-exp-nested-subalg*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f \text{subalgebra } M G \text{subalgebra } G F$
shows $AE \xi \text{ in } M. \text{cond-exp } M F f \xi = \text{cond-exp } M F (\text{cond-exp } M G f) \xi$
using *has-cond-expI assms sigma-finite-subalgebra-def* **by** (*auto intro!*: *has-cond-exp-nested-subalg* [*THEN has-cond-exp-charact(2), THEN AE-symmetric*] *sigma-finite-subalgebra.has-cond-expI* [*OF sigma-finite-subalgebra.intro* [*OF assms(2)*]] *nested-subalg-is-sigma-finite*)

The conditional expectation is linear.

lemma *cond-exp-add*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f \text{integrable } M g$
shows $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x + g x) x = \text{cond-exp } M F f x + \text{cond-exp } M F g x$
using *has-cond-exp-add* [*OF has-cond-expI(1,1), OF assms, THEN has-cond-exp-charact(2)*]
.

lemma *cond-exp-diff*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f \text{integrable } M g$
shows $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x - g x) x = \text{cond-exp } M F f x - \text{cond-exp } M F g x$
using *has-cond-exp-add* [*OF - has-cond-exp-scaleR-right, OF has-cond-expI(1,1), OF assms, THEN has-cond-exp-charact(2), of -1*] **by** *simp*

lemma *cond-exp-diff'*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f \text{integrable } M g$
shows $AE x \text{ in } M. \text{cond-exp } M F (f - g) x = \text{cond-exp } M F f x - \text{cond-exp } M F g x$
unfolding *fun-diff-def* **using** *assms* **by** (*rule cond-exp-diff*)

lemma *cond-exp-scaleR-left*:
fixes $f :: 'a \Rightarrow \text{real}$
assumes *integrable* $M f$
shows $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x *_{\mathbb{R}} c) x = \text{cond-exp } M F f x *_{\mathbb{R}} c$

using *cond-exp-set-integral*[*OF assms*] *subalg assms* **unfolding** *subalgebra-def*
by (*intro cond-exp-charact*,
subst set-integral-scaleR-left, *blast*, *intro assms*,
subst set-integral-scaleR-left, *blast*, *intro integrable-cond-exp*)
auto

The conditional expectation operator is a contraction, i.e. a bounded linear operator with operator norm less than or equal to 1.

To show this we first obtain a subsequence $\lambda x i. s (r i) x$, such that $\lambda i. \text{cond-exp } M F (s (r i)) x$ converges to $\text{cond-exp } M F f x$ a.e. Afterwards, we obtain a sub-subsequence $\lambda x i. s (r (r' i)) x$, such that $\lambda i. \text{cond-exp } M F (\lambda x. \text{norm } (s (r i))) x$ converges to $\text{cond-exp } M F (\lambda x. \text{norm } (f x)) x$ a.e. Finally, we show that the inequality holds by showing that the terms of the subsequences obey the inequality and the fact that a subsequence of a convergent sequence converges to the same limit.

lemma *cond-exp-contraction*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$

assumes *integrable M f*

shows $AE x \text{ in } M. \text{norm } (\text{cond-exp } M F f x) \leq \text{cond-exp } M F (\lambda x. \text{norm } (f x))$

x

proof –

obtain s **where** $s: \bigwedge i. \text{simple-function } M (s i) \bigwedge i. \text{emeasure } M \{y \in \text{space } M. s i y \neq 0\} \neq \infty \bigwedge x. x \in \text{space } M \implies (\lambda i. s i x) \longrightarrow f x \bigwedge i x. x \in \text{space } M \implies \text{norm } (s i x) \leq 2 * \text{norm } (f x)$

by (*blast intro: integrable-implies-simple-function-sequence*[*OF assms*])

obtain r **where** $r: \text{strict-mono } r$ **and** $\text{has-cond-exp } M F f (\lambda x. \lim (\lambda i. \text{cond-exp } M F (s (r i)) x)) AE x \text{ in } M. (\lambda i. \text{cond-exp } M F (s (r i)) x) \longrightarrow \lim (\lambda i. \text{cond-exp } M F (s (r i)) x)$

using *has-cond-exp-simple-lim*[*OF assms s*] **unfolding** *convergent-LIMSEQ-iff*
by *blast*

hence $r\text{-tendsto}: AE x \text{ in } M. (\lambda i. \text{cond-exp } M F (s (r i)) x) \longrightarrow \text{cond-exp } M F f x$ **using** *has-cond-exp-charact(2)* **by** *force*

have $\text{norm-s-r}: \bigwedge i. \text{simple-function } M (\lambda x. \text{norm } (s (r i) x)) \bigwedge i. \text{emeasure } M \{y \in \text{space } M. \text{norm } (s (r i) y) \neq 0\} \neq \infty \bigwedge x. x \in \text{space } M \implies (\lambda i. \text{norm } (s (r i) x)) \longrightarrow \text{norm } (f x) \bigwedge i x. x \in \text{space } M \implies \text{norm } (\text{norm } (s (r i) x)) \leq 2 * \text{norm } (\text{norm } (f x))$

using s **by** (*auto intro: LIMSEQ-subseq-LIMSEQ*[*OF tendsto-norm r, unfolded comp-def*] *simple-function-compose1*)

obtain r' **where** $r': \text{strict-mono } r'$ **and** $\text{has-cond-exp } M F (\lambda x. \text{norm } (f x)) (\lambda x. \lim (\lambda i. \text{cond-exp } M F (\lambda x. \text{norm } (s (r (r' i)) x)) x)) AE x \text{ in } M. (\lambda i. \text{cond-exp } M F (\lambda x. \text{norm } (s (r (r' i)) x)) x) \longrightarrow \lim (\lambda i. \text{cond-exp } M F (\lambda x. \text{norm } (s (r (r' i)) x)) x)$ **using** *has-cond-exp-simple-lim*[*OF integrable-norm norm-s-r, OF assms*] **unfolding** *convergent-LIMSEQ-iff* **by** *blast*

hence $r'\text{-tendsto}: AE x \text{ in } M. (\lambda i. \text{cond-exp } M F (\lambda x. \text{norm } (s (r (r' i)) x)) x) \longrightarrow \text{cond-exp } M F (\lambda x. \text{norm } (f x)) x$ **using** *has-cond-exp-charact(2)* **by** *force*

have $AE\ x\ in\ M. \forall i. norm\ (cond-exp\ M\ F\ (s\ (r\ (r'\ i)))\ x) \leq cond-exp\ M\ F\ (\lambda x. norm\ (s\ (r\ (r'\ i)))\ x)\ x$ **using** s **by** $(auto\ intro: cond-exp-contraction-simple\ simp\ add: AE-all-countable)$
moreover have $AE\ x\ in\ M. (\lambda i. norm\ (cond-exp\ M\ F\ (s\ (r\ (r'\ i)))\ x)) \longrightarrow norm\ (cond-exp\ M\ F\ f\ x)$ **using** $r-tendsto\ LIMSEQ-subseq-LIMSEQ[OF\ tendsto-norm\ r',\ unfolded\ comp-def]$ **by** $fast$
ultimately show $?thesis$ **using** $LIMSEQ-le\ r'-tendsto$ **by** $fast$
qed

The following lemmas are called "pulling out whats known". We first show the statement for real-valued functions using the lemma *real-cond-exp-intg*, which is already present. We then show it for arbitrary g using the lecture notes of Gordan Zitkovic for the course "Theory of Probability I" \cite{Zitkovic-2015}.

lemma *cond-exp-measurable-mult*:

fixes $f\ g :: 'a \Rightarrow real$

assumes $[measurable]: integrable\ M\ (\lambda x. f\ x * g\ x)\ integrable\ M\ g\ f \in borel-measurable\ F$

shows $integrable\ M\ (\lambda x. f\ x * cond-exp\ M\ F\ g\ x)$

$AE\ x\ in\ M. cond-exp\ M\ F\ (\lambda x. f\ x * g\ x)\ x = f\ x * cond-exp\ M\ F\ g\ x$

proof –

show $integrable: integrable\ M\ (\lambda x. f\ x * cond-exp\ M\ F\ g\ x)$ **using** $cond-exp-real[OF\ assms(2)]$ **by** $(intro\ integrable-cong-AE-imp[OF\ real-cond-exp-intg(1),\ OF\ assms(1,3)]\ assms(2)[THEN\ borel-measurable-integrable])\ measurable-from-subalg[OF\ subalg])\ auto$

interpret $sigma-finite-measure\ restr-to-subalg\ M\ F$ **by** $(rule\ sigma-fin-subalg)$

{

fix A **assume** $asm: A \in sets\ F$

hence $asm': A \in sets\ M$ **using** $subalg$ **by** $(fastforce\ simp\ add: subalgebra-def)$

have $set-lebesgue-integral\ M\ A\ (cond-exp\ M\ F\ (\lambda x. f\ x * g\ x)) = set-lebesgue-integral\ M\ A\ (\lambda x. f\ x * g\ x)$ **by** $(simp\ add: cond-exp-set-integral[OF\ assms(1)\ asm])$

also have $\dots = set-lebesgue-integral\ M\ A\ (\lambda x. f\ x * real-cond-exp\ M\ F\ g\ x)$ **using** $borel-measurable-times[OF\ borel-measurable-indicator[OF\ asm]\ assms(3)]\ borel-measurable-integrable[OF\ assms(2)]\ integrable-mult-indicator[OF\ asm'\ assms(1)]$

by $(fastforce\ simp\ add: set-lebesgue-integral-def\ mult.assoc[symmetric]\ intro: real-cond-exp-intg(2)[symmetric])$

also have $\dots = set-lebesgue-integral\ M\ A\ (\lambda x. f\ x * cond-exp\ M\ F\ g\ x)$ **using** $cond-exp-real[OF\ assms(2)]\ asm'\ borel-measurable-cond-exp'\ borel-measurable-cond-exp2\ measurable-from-subalg[OF\ subalg\ assms(3)]$ **by** $(auto\ simp\ add: set-lebesgue-integral-def\ intro: integral-cong-AE)$

finally have $set-lebesgue-integral\ M\ A\ (cond-exp\ M\ F\ (\lambda x. f\ x * g\ x)) = \int_{x \in A}. (f\ x * cond-exp\ M\ F\ g\ x) dM$.

}

hence $AE\ x\ in\ restr-to-subalg\ M\ F. cond-exp\ M\ F\ (\lambda x. f\ x * g\ x)\ x = f\ x * cond-exp\ M\ F\ g\ x$ **by** $(intro\ density-unique-banach\ integrable-cond-exp\ integrable\ integrable-in-subalg\ subalg,\ measurable,\ simp\ add: set-lebesgue-integral-def\ integral-subalgebra2[OF\ subalg]\ sets-restr-to-subalg[OF\ subalg])$

thus $AE\ x\ in\ M. cond-exp\ M\ F\ (\lambda x. f\ x * g\ x)\ x = f\ x * cond-exp\ M\ F\ g\ x$ **by** $(rule\ AE-restr-to-subalg[OF\ subalg])$

qed

lemma *cond-exp-measurable-scaleR*:

fixes $f :: 'a \Rightarrow \text{real}$ **and** $g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes $[\text{measurable}]$: $\text{integrable } M (\lambda x. f x *_R g x)$ $\text{integrable } M g$ $f \in \text{borel-measurable}$
 F

shows $\text{integrable } M (\lambda x. f x *_R \text{cond-exp } M F g x)$

$AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x *_R g x) x = f x *_R \text{cond-exp } M F g x$

proof –

let $?F = \text{restr-to-subalg } M F$

have subalg' : $\text{subalgebra } M (\text{restr-to-subalg } M F)$ **by** $(\text{metis sets-eq-imp-space-eq sets-restr-to-subalg subalg subalgebra-def})$

{

fix z **assume** $\text{asm}[\text{measurable}]$: $\text{integrable } M (\lambda x. z x *_R g x)$ $z \in \text{borel-measurable}$
 $?F$

hence $\text{asm}'[\text{measurable}]$: $z \in \text{borel-measurable } F$ **using** $\text{measurable-in-subalg}'$
 subalg **by** blast

have $\text{integrable } M (\lambda x. z x *_R \text{cond-exp } M F g x)$ $LINT x | M. z x *_R g x =$
 $LINT x | M. z x *_R \text{cond-exp } M F g x$

proof –

obtain s **where** $s\text{-is}$: $\bigwedge i. \text{simple-function } ?F (s i) \bigwedge x. x \in \text{space } ?F \implies (\lambda i. s i x) \longrightarrow z x \bigwedge i x. x \in \text{space } ?F \implies \text{norm } (s i x) \leq 2 * \text{norm } (z x)$ **using**
 $\text{borel-measurable-implies-sequence-metric}[OF \text{asm}(2), of 0]$ **by** force

We need to apply the dominated convergence theorem twice, therefore we need to show the following prerequisites.

have $s\text{-scaleR-g-tendsto}$: $AE x \text{ in } M. (\lambda i. s i x *_R g x) \longrightarrow z x *_R g x$
using $s\text{-is}(2)$ **by** $(\text{simp add: space-restr-to-subalg tendsto-scaleR})$

have $s\text{-scaleR-cond-exp-g-tendsto}$: $AE x \text{ in } ?F. (\lambda i. s i x *_R \text{cond-exp } M F g x) \longrightarrow z x *_R \text{cond-exp } M F g x$ **using** $s\text{-is}(2)$ **by** $(\text{simp add: tendsto-scaleR})$

have $s\text{-scaleR-g-meas}$: $(\lambda x. s i x *_R g x) \in \text{borel-measurable } M$ **for** i **using** $s\text{-is}(1)[THEN \text{borel-measurable-simple-function}, THEN \text{subalg}'[THEN \text{measurable-from-subalg}]]$ **by** simp

have $s\text{-scaleR-cond-exp-g-meas}$: $(\lambda x. s i x *_R \text{cond-exp } M F g x) \in \text{borel-measurable}$
 $?F$ **for** i **using** $s\text{-is}(1)[THEN \text{borel-measurable-simple-function}]$ $\text{measurable-in-subalg}[OF \text{subalg borel-measurable-cond-exp}]$ **by** $(\text{fastforce intro: borel-measurable-scaleR})$

have $s\text{-scaleR-g-AE-bdd}$: $AE x \text{ in } M. \text{norm } (s i x *_R g x) \leq 2 * \text{norm } (z x *_R g x)$ **for** i **using** $s\text{-is}(3)$ **by** $(\text{fastforce simp add: space-restr-to-subalg mult.assoc[symmetric] mult-right-mono})$

{

fix i

have asm : $\text{integrable } M (\lambda x. \text{norm } (z x) * \text{norm } (g x))$ **using** $\text{asm}(1)[THEN \text{integrable-norm}]$ **by** simp

have $AE x \text{ in } ?F. \text{norm } (s i x *_R \text{cond-exp } M F g x) \leq 2 * \text{norm } (z x) * \text{norm } (\text{cond-exp } M F g x)$ **using** $s\text{-is}(3)$ **by** $(\text{fastforce simp add: mult-mono})$

moreover have $AE x \text{ in } ?F. \text{norm } (z x) * \text{cond-exp } M F (\lambda x. \text{norm } (g x)) x = \text{cond-exp } M F (\lambda x. \text{norm } (z x) * \text{norm } (g x)) x$ **by** $(\text{rule cond-exp-measurable-mult}(2)[THEN$

AE-symmetric, OF asm integrable-norm, OF assms(2), THEN AE-restr-to-subalg2[OF subalg]], auto)

ultimately have *AE x in ?F. norm (s i x *_R cond-exp M F g x) ≤ 2 * cond-exp M F (λx. norm (z x *_R g x)) x using cond-exp-contraction[OF assms(2), THEN AE-restr-to-subalg2[OF subalg]] order-trans[OF - mult-mono] by fastforce*
}
note *s-scaleR-cond-exp-g-AE-bdd = this*

In the following section we need to pay attention to which measures we are using for integration. The rhs is F-measurable while the lhs is only M-measurable.

{
fix *i*
have *s-meas-M[measurable]: s i ∈ borel-measurable M by (meson borel-measurable-simple-function measurable-from-subalg s-is(1) subalg')*
have *s-meas-F[measurable]: s i ∈ borel-measurable F by (meson borel-measurable-simple-function measurable-in-subalg' s-is(1) subalg)*

have *s-scaleR-eq: s i x *_R h x = (∑ y ∈ s i ' space M. (indicator (s i - ' {y} ∩ space M) x *_R y) *_R h x) if x ∈ space M for x and h :: 'a ⇒ 'b using simple-function-indicator-representation[OF s-is(1), of x i] that unfolding space-restr-to-subalg scaleR-left.sum[of - - h x, symmetric] by presburger*

have *LINT x|M. s i x *_R g x = LINT x|M. (∑ y ∈ s i ' space M. indicator (s i - ' {y} ∩ space M) x *_R y *_R g x) using s-scaleR-eq by (intro Bochner-Integration.integral-cong) auto*
also have *... = (∑ y ∈ s i ' space M. LINT x|M. indicator (s i - ' {y} ∩ space M) x *_R y *_R g x) by (intro Bochner-Integration.integral-sum integrable-mult-indicator[OF - integrable-scaleR-right] assms(2)) simp*
also have *... = (∑ y ∈ s i ' space M. y *_R set-lebesgue-integral M (s i - ' {y} ∩ space M) g) by (simp only: set-lebesgue-integral-def[symmetric]) simp*
also have *... = (∑ y ∈ s i ' space M. y *_R set-lebesgue-integral M (s i - ' {y} ∩ space M) (cond-exp M F g)) using assms(2) subalg borel-measurable-vimage[OF s-meas-F] by (subst cond-exp-set-integral, auto simp add: subalgebra-def)*
also have *... = (∑ y ∈ s i ' space M. LINT x|M. indicator (s i - ' {y} ∩ space M) x *_R y *_R cond-exp M F g x) by (simp only: set-lebesgue-integral-def[symmetric]) simp*
also have *... = LINT x|M. (∑ y ∈ s i ' space M. indicator (s i - ' {y} ∩ space M) x *_R y *_R cond-exp M F g x) by (intro Bochner-Integration.integral-sum[symmetric] integrable-mult-indicator[OF - integrable-scaleR-right]) auto*
also have *... = LINT x|M. s i x *_R cond-exp M F g x using s-scaleR-eq by (intro Bochner-Integration.integral-cong) auto*
finally have *LINT x|M. s i x *_R g x = LINT x|?F. s i x *_R cond-exp M F g x by (simp add: integral-subalgebra2[OF subalg])*
}
note *integral-s-eq = this*

Now we just plug in the results we obtained into DCT, and use the fact that limits are unique.

show *integrable* M $(\lambda x. z x *_R \text{cond-exp } M F g x)$ **using** *s-scaleR-cond-exp-g-meas* *asm*(2) *borel-measurable-cond-exp'* **by** (intro *integrable-from-subalg*[*OF subalg*] *integrable-cond-exp* *integrable-dominated-convergence*[*OF - - s-scaleR-cond-exp-g-tendsto* *s-scaleR-cond-exp-g-AE-bdd*]) (auto intro: *measurable-from-subalg*[*OF subalg*] *integrable-in-subalg* *measurable-in-subalg* *subalg*)

have $(\lambda i. LINT x|M. s i x *_R g x) \longrightarrow LINT x|M. z x *_R g x$ **using** *s-scaleR-g-meas* *asm*(1)[*THEN integrable-norm*] *asm'* *borel-measurable-cond-exp'* **by** (intro *integral-dominated-convergence*[*OF - - s-scaleR-g-tendsto* *s-scaleR-g-AE-bdd*]) (auto intro: *measurable-from-subalg*[*OF subalg*])

moreover have $(\lambda i. LINT x|?F. s i x *_R \text{cond-exp } M F g x) \longrightarrow LINT x|?F. z x *_R \text{cond-exp } M F g x$ **using** *s-scaleR-cond-exp-g-meas* *asm*(2) *borel-measurable-cond-exp'* **by** (intro *integral-dominated-convergence*[*OF - - s-scaleR-cond-exp-g-tendsto* *s-scaleR-cond-exp-g-AE-bdd*]) (auto intro: *measurable-from-subalg*[*OF subalg*] *integrable-in-subalg* *measurable-in-subalg* *subalg*)

ultimately show $LINT x|M. z x *_R g x = LINT x|M. z x *_R \text{cond-exp } M F g x$ **using** *integral-s-eq* **using** *subalg* **by** (*simp add: LIMSEQ-unique integral-subalgebra2*)

qed

}

note $*$ = *this*

The main statement now follows with $z = (\lambda x. \text{indicat-real } A x * f x)$.

show *integrable* M $(\lambda x. f x *_R \text{cond-exp } M F g x)$ **using** $*$ *assms* *measurable-in-subalg*[*OF subalg*] **by** *blast*

{

fix A **assume** *asm*: $A \in F$

hence *integrable* M $(\lambda x. \text{indicat-real } A x *_R f x *_R g x)$ **using** *subalg* **by** (*fastforce simp add: subalgebra-def intro!: integrable-mult-indicator assms*(1))

hence *set-lebesgue-integral* $M A$ $(\lambda x. f x *_R g x) = \text{set-lebesgue-integral } M A$ $(\lambda x. f x *_R \text{cond-exp } M F g x)$ **unfolding** *set-lebesgue-integral-def* **using** *asm* **by** (auto intro!: $*$ *measurable-in-subalg*[*OF subalg*])

}

thus *AE* x *in* $M. \text{cond-exp } M F$ $(\lambda x. f x *_R g x) x = f x *_R \text{cond-exp } M F g x$ **using** *borel-measurable-cond-exp* **by** (intro *cond-exp-charact*, auto intro!: $*$ *assms* *measurable-in-subalg*[*OF subalg*])

qed

lemma *cond-exp-sum* [intro, simp]:

fixes $f :: 't \Rightarrow 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$

assumes [*measurable*]: $\bigwedge i. \text{integrable } M (f i)$

shows *AE* x *in* $M. \text{cond-exp } M F$ $(\lambda x. \sum_{i \in I. f i x}) x = (\sum_{i \in I. \text{cond-exp } M F (f i) x})$

proof (*rule has-cond-exp-charact*, intro *has-cond-expI'*)

fix A **assume** [*measurable*]: $A \in \text{sets } F$

then have *A-meas* [*measurable*]: $A \in \text{sets } M$ **by** (*meson subsetD subalg subalgebra-def*)

have $(\int x \in A. (\sum i \in I. f i x) \partial M) = (\int x. (\sum i \in I. \text{indicator } A x *_R f i x) \partial M)$
unfolding *set-lebesgue-integral-def* **by** (*simp add: scaleR-sum-right*)
also have $\dots = (\sum i \in I. (\int x. \text{indicator } A x *_R f i x \partial M))$ **using** *assms* **by** (*auto intro!: Bochner-Integration.integral-sum integrable-mult-indicator*)
also have $\dots = (\sum i \in I. (\int x. \text{indicator } A x *_R \text{cond-exp } M F (f i) x \partial M))$ **using** *cond-exp-set-integral[OF assms]* **by** (*simp add: set-lebesgue-integral-def*)
also have $\dots = (\int x. (\sum i \in I. \text{indicator } A x *_R \text{cond-exp } M F (f i) x) \partial M)$
using *assms* **by** (*auto intro!: Bochner-Integration.integral-sum[symmetric] integrable-mult-indicator*)
also have $\dots = (\int x \in A. (\sum i \in I. \text{cond-exp } M F (f i) x) \partial M)$ **unfolding** *set-lebesgue-integral-def*
by (*simp add: scaleR-sum-right*)
finally show $(\int x \in A. (\sum i \in I. f i x) \partial M) = (\int x \in A. (\sum i \in I. \text{cond-exp } M F (f i) x) \partial M)$ **by** *auto*
qed (*auto simp add: assms integrable-cond-exp*)

4.3 Linearly Ordered Banach Spaces

In this subsection we show monotonicity results concerning the conditional expectation operator.

lemma *cond-exp-gr-c*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes *integrable M f AE x in M. f x > c*

shows *AE x in M. cond-exp M F f x > c*

proof –

define X **where** $X = \{x \in \text{space } M. \text{cond-exp } M F f x \leq c\}$

have [*measurable*]: $X \in \text{sets } F$ **unfolding** *X-def* **by** *measurable (metis sets.top subalg subalgebra-def)*

hence $X \text{-in-} M$: $X \in \text{sets } M$ **using** *sets-restr-to-subalg subalg subalgebra-def* **by** *blast*

have *emeasure M X = 0*

proof (*rule ccontr*)

assume *emeasure M X ≠ 0*

have *emeasure (restr-to-subalg M F) X = emeasure M X* **by** (*simp add: emeasure-restr-to-subalg subalg*)

hence *emeasure (restr-to-subalg M F) X > 0* **using** $\neg(\text{emeasure } M X = 0)$ **by** *gr-zeroI* **by** *auto*

then obtain A **where** $A \in \text{sets } (restr\text{-to-subalg } M F)$ $A \subseteq X$ *emeasure (restr-to-subalg M F) A > 0* *emeasure (restr-to-subalg M F) A < ∞*

using *sigma-fin-subalg* **by** (*metis emeasure-notin-sets ennreal-0 infinity-ennreal-def le-less-linear neq-top-trans not-gr-zero order-refl sigma-finite-measure.approx-PInf-emeasure-with-finite*)

hence [*simp*]: $A \in \text{sets } F$ **using** *subalg sets-restr-to-subalg* **by** *blast*

hence $A \text{-in-sets-} M$ [*simp*]: $A \in \text{sets } M$ **using** *sets-restr-to-subalg subalg subalgebra-def* **by** *blast*

have [*simp*]: *set-integrable M A* $(\lambda x. c)$ **using** *A subalg* **by** (*auto simp add: set-integrable-def emeasure-restr-to-subalg*)

have [*simp*]: *set-integrable M A f* **unfolding** *set-integrable-def* **by** (*rule integrable-mult-indicator, auto simp add: assms(1)*)

have *AE x in M. indicator A x *_R c = indicator A x *_R f x*

proof (rule integral-eq-mono-AE-eq-AE)
have $(\int x \in A. c \partial M) \leq (\int x \in A. f x \partial M)$ **using** *assms(2)* **by** (intro set-integral-mono-AE-banach)
auto
moreover
{
have $(\int x \in A. f x \partial M) = (\int x \in A. \text{cond-exp } M F f x \partial M)$ **by** (rule
cond-exp-set-integral, auto simp add: assms)
also have $\dots \leq (\int x \in A. c \partial M)$ **using** *A* **by** (auto intro!: set-integral-mono-banach
simp add: X-def)
finally have $(\int x \in A. f x \partial M) \leq (\int x \in A. c \partial M)$ **by** *simp*
}
ultimately show $LINT x | M. \text{indicator } A x *_R c = LINT x | M. \text{indicator } A$
 $x *_R f x$ **unfolding** *set-lebesgue-integral-def* **by** *simp*
show $AE x \text{ in } M. \text{indicator } A x *_R c \leq \text{indicator } A x *_R f x$ **using** *assms* **by**
(auto simp add: X-def indicator-def)
qed (auto simp add: set-integrable-def[symmetric])
hence $AE x \in A \text{ in } M. c = f x$ **by** *auto*
hence $AE x \in A \text{ in } M. \text{False}$ **using** *assms(2)* **by** *auto*
hence $A \in \text{null-sets } M$ **using** *AE-iff-null-sets A-in-sets-M* **by** *metis*
thus *False* **using** *A(3)* **by** (simp add: emeasure-restr-to-subalg null-setsD1
subalg)
qed
thus *?thesis* **using** *AE-iff-null-sets[OF X-in-M]* **unfolding** *X-def* **by** *auto*
qed

corollary *cond-exp-less-c*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$
dered-real-vector}
assumes *integrable M f AE x in M. f x < c*
shows $AE x \text{ in } M. \text{cond-exp } M F f x < c$
proof –
have $AE x \text{ in } M. \text{cond-exp } M F f x = - \text{cond-exp } M F (\lambda x. - f x) x$ **using**
cond-exp-uminus[OF assms(1)] **by** *auto*
moreover have $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. - f x) x > - c$ **using** *assms*
by (intro *cond-exp-gr-c*) *auto*
ultimately show *?thesis* **by** (force simp add: minus-less-iff)
qed

lemma *cond-exp-mono-strict*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$
dered-real-vector}
assumes *integrable M f integrable M g AE x in M. f x < g x*
shows $AE x \text{ in } M. \text{cond-exp } M F f x < \text{cond-exp } M F g x$
using *cond-exp-less-c[OF Bochner-Integration.integrable-diff, OF assms(1,2), of*
0]
cond-exp-diff[OF assms(1,2)] assms(3) **by** *auto*

lemma *cond-exp-ge-c*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$

dered-real-vector }
assumes $[measurable]: \text{integrable } M f$
and $AE\ x\ in\ M. f\ x \geq c$
shows $AE\ x\ in\ M. \text{cond-exp } M\ F\ f\ x \geq c$
proof –
let $?F = \text{restr-to-subalg } M\ F$
interpret $\text{sigma-finite-measure restr-to-subalg } M\ F$ **using** sigma-fin-subalg **by**
 auto
{
fix A **assume** $\text{asm}: A \in \text{sets } ?F\ 0 < \text{measure } ?F\ A$
have $[simp]: \text{sets } ?F = \text{sets } F\ \text{measure } ?F\ A = \text{measure } M\ A$ **using** asm **by** $(\text{auto } \text{simp add: measure-def sets-restr-to-subalg}[OF\ \text{subalg}] \text{emeasure-restr-to-subalg}[OF\ \text{subalg}])$
have $M-A: \text{emeasure } M\ A < \infty$ **using** $\text{measure-zero-top asm}$ **by** $(\text{force simp add: top.not-eq-extremum})$
hence $F-A: \text{emeasure } ?F\ A < \infty$ **using** $\text{asm}(1) \text{emeasure-restr-to-subalg subalg}$
by fastforce
have $\text{set-lebesgue-integral } M\ A\ (\lambda\cdot. c) \leq \text{set-lebesgue-integral } M\ A\ f$ **using**
 $\text{assms asm } M-A\ \text{subalg}$ **by** $(\text{intro set-integral-mono-AE-banach, auto simp add: set-integrable-def integrable-mult-indicator subalgebra-def sets-restr-to-subalg})$
also have $\dots = \text{set-lebesgue-integral } M\ A\ (\text{cond-exp } M\ F\ f)$ **using** $\text{cond-exp-set-integral}[OF\ \text{assms}(1)]\ \text{asm}$ **by** auto
also have $\dots = \text{set-lebesgue-integral } ?F\ A\ (\text{cond-exp } M\ F\ f)$ **unfolding** $\text{set-lebesgue-integral-def}$
using $\text{asm borel-measurable-cond-exp}$ **by** $(\text{intro integral-subalgebra2}[OF\ \text{subalg, symmetric}], \text{simp})$
finally have $(1 / \text{measure } ?F\ A) *_R \text{set-lebesgue-integral } ?F\ A\ (\text{cond-exp } M\ F\ f)$
 $\in \{c..\}$ **using** $\text{asm subalg } M-A$ **by** $(\text{auto simp add: set-integral-const subalgebra-def intro!: pos-divideR-le-eq}[THEN iffD1])$
}
thus $?thesis$ **using** $\text{AE-restr-to-subalg}[OF\ \text{subalg}] \text{averaging-theorem}[OF\ \text{integrable-in-subalg closed-atLeast, OF subalg borel-measurable-cond-exp integrable-cond-exp}]$
by auto
qed

corollary cond-exp-le-c :

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes $\text{integrable } M\ f$
and $AE\ x\ in\ M. f\ x \leq c$
shows $AE\ x\ in\ M. \text{cond-exp } M\ F\ f\ x \leq c$
proof –
have $AE\ x\ in\ M. \text{cond-exp } M\ F\ f\ x = - \text{cond-exp } M\ F\ (\lambda x. - f\ x)\ x$ **using**
 $\text{cond-exp-uminus}[OF\ \text{assms}(1)]$ **by** force
moreover have $AE\ x\ in\ M. \text{cond-exp } M\ F\ (\lambda x. - f\ x)\ x \geq - c$ **using** assms
by $(\text{intro cond-exp-ge-c})\ \text{auto}$
ultimately show $?thesis$ **by** $(\text{force simp add: minus-le-iff})$
qed

corollary cond-exp-mono :

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes $\text{integrable } M f \text{ integrable } M g \text{ } AE \ x \text{ in } M. f \ x \leq g \ x$
shows $AE \ x \text{ in } M. \text{cond-exp } M F f \ x \leq \text{cond-exp } M F g \ x$
using $\text{cond-exp-le-c}[OF \ \text{Bochner-Integration.integrable-diff}, OF \ \text{assms}(1,2), \text{ of } 0]$
 $\text{cond-exp-diff}[OF \ \text{assms}(1,2)] \ \text{assms}(3) \text{ by auto}$

corollary *cond-exp-min:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes $\text{integrable } M f \text{ integrable } M g$
shows $AE \ \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \min (f \ x) (g \ x)) \ \xi \leq \min (\text{cond-exp } M F f \ \xi) (\text{cond-exp } M F g \ \xi)$
proof –
have $AE \ \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \min (f \ x) (g \ x)) \ \xi \leq \text{cond-exp } M F f \ \xi \text{ by } (intro \ \text{cond-exp-mono integrable-min assms, simp})$
moreover have $AE \ \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \min (f \ x) (g \ x)) \ \xi \leq \text{cond-exp } M F g \ \xi \text{ by } (intro \ \text{cond-exp-mono integrable-min assms, simp})$
ultimately show $AE \ \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \min (f \ x) (g \ x)) \ \xi \leq \min (\text{cond-exp } M F f \ \xi) (\text{cond-exp } M F g \ \xi) \text{ by fastforce}$
qed

corollary *cond-exp-max:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes $\text{integrable } M f \text{ integrable } M g$
shows $AE \ \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \max (f \ x) (g \ x)) \ \xi \geq \max (\text{cond-exp } M F f \ \xi) (\text{cond-exp } M F g \ \xi)$
proof –
have $AE \ \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \max (f \ x) (g \ x)) \ \xi \geq \text{cond-exp } M F f \ \xi \text{ by } (intro \ \text{cond-exp-mono integrable-max assms, simp})$
moreover have $AE \ \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \max (f \ x) (g \ x)) \ \xi \geq \text{cond-exp } M F g \ \xi \text{ by } (intro \ \text{cond-exp-mono integrable-max assms, simp})$
ultimately show $AE \ \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \max (f \ x) (g \ x)) \ \xi \geq \max (\text{cond-exp } M F f \ \xi) (\text{cond-exp } M F g \ \xi) \text{ by fastforce}$
qed

corollary *cond-exp-inf:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector, lattice}\}$
assumes $\text{integrable } M f \text{ integrable } M g$
shows $AE \ \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \inf (f \ x) (g \ x)) \ \xi \leq \inf (\text{cond-exp } M F f \ \xi) (\text{cond-exp } M F g \ \xi)$
unfolding *inf-min* **using** *assms* **by** (*rule cond-exp-min*)

corollary *cond-exp-sup:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector, lattice}\}$

```

    assumes integrable M f integrable M g
    shows AE  $\xi$  in M. cond-exp M F ( $\lambda x. \sup (f x) (g x)$ )  $\xi \geq \sup (cond-exp M F f$ 
 $\xi) (cond-exp M F g \xi)$ 
    unfolding sup-max using assms by (rule cond-exp-max)

end

```

4.4 Probability Spaces

```

lemma (in prob-space) sigma-finite-subalgebra-restr-to-subalg:
  assumes subalgebra M F
  shows sigma-finite-subalgebra M F
proof (intro sigma-finite-subalgebra.intro)
  interpret F: prob-space restr-to-subalg M F using assms prob-space-restr-to-subalg
  prob-space-axioms by blast
  show sigma-finite-measure (restr-to-subalg M F) by (rule F.sigma-finite-measure-axioms)
qed (rule assms)

```

```

lemma (in prob-space) cond-exp-trivial:
  fixes f :: 'a  $\Rightarrow$  'b :: {second-countable-topology, banach}
  assumes integrable M f
  shows AE x in M. cond-exp M (sigma (space M) {}) f x = expectation f
proof -
  interpret sigma-finite-subalgebra M sigma (space M) {} by (auto intro: sigma-finite-subalgebra-restr-to-subalg)
  simp add: subalgebra-def sigma-sets-empty-eq
  show ?thesis using assms by (intro cond-exp-charact) (auto simp add: sigma-sets-empty-eq
  set-lebesgue-integral-def prob-space cong: Bochner-Integration.integral-cong)
qed

```

The following lemma shows that independent σ -algebras don't matter for the conditional expectation. The proof is adapted from \cite{Zitkovic-2015}.

```

lemma (in prob-space) cond-exp-indep-subalgebra:
  fixes f :: 'a  $\Rightarrow$  'b :: {second-countable-topology, banach, real-normed-field}
  assumes subalgebra: subalgebra M F subalgebra M G
    and independent: indep-set G (sigma (space M) (F  $\cup$  vimage-algebra (space
M) f borel))
  assumes [measurable]: integrable M f
  shows AE x in M. cond-exp M (sigma (space M) (F  $\cup$  G)) f x = cond-exp M F
f x
proof -
  interpret Un-sigma: sigma-finite-subalgebra M sigma (space M) (F  $\cup$  G) using
  assms(1,2) by (auto intro!: sigma-finite-subalgebra-restr-to-subalg sets.sigma-sets-subset
  simp add: subalgebra-def space-measure-of-conv sets-measure-of-conv)
  interpret sigma-finite-subalgebra M F using assms by (auto intro: sigma-finite-subalgebra-restr-to-subalg)
  {
    fix A
    assume asm: A  $\in$  sigma (space M) {a  $\cap$  b | a b. a  $\in$  F  $\wedge$  b  $\in$  G}
    have in-events: sigma-sets (space M) {a  $\cap$  b | a b. a  $\in$  sets F  $\wedge$  b  $\in$  sets
    G}  $\subseteq$  events using subalgebra by (intro sets.sigma-sets-subset, auto simp add:

```

```

subalgebra-def)
  have Int-stable  $\{a \cap b \mid a \text{ b. } a \in F \wedge b \in G\}$ 
  proof -
    {
      fix af bf ag bg
      assume F: af  $\in F$  bf  $\in F$  and G: ag  $\in G$  bg  $\in G$ 
      have af  $\cap$  bf  $\in F$  by (intro sets.Int F)
      moreover have ag  $\cap$  bg  $\in G$  by (intro sets.Int G)
      ultimately have  $\exists a \text{ b. } af \cap ag \cap (bf \cap bg) = a \cap b \wedge a \in \text{sets } F \wedge b \in \text{sets } G$  by (metis inf-assoc inf-left-commute)
    }
    thus ?thesis by (force intro!: Int-stableI)
  qed
  moreover have  $\{a \cap b \mid a \text{ b. } a \in F \wedge b \in G\} \subseteq \text{Pow } (\text{space } M)$  using
  subalgebra by (force simp add: subalgebra-def dest: sets.sets-into-space)
  moreover have  $A \in \text{sigma-sets } (\text{space } M) \{a \cap b \mid a \text{ b. } a \in F \wedge b \in G\}$  using
  calculation asm by force
  ultimately have  $\text{set-lebesgue-integral } M \ A \ f = \text{set-lebesgue-integral } M \ A$ 
  (cond-exp  $M \ F \ f$ )
  proof (induction rule: sigma-sets-induct-disjoint)
    case (basic A)
    then obtain a b where A:  $A = a \cap b \ a \in F \ b \in G$  by blast

    hence events[measurable]:  $a \in \text{events } b \in \text{events}$  using subalgebra by (auto
    simp add: subalgebra-def)

    have [simp]:  $\text{sigma-sets } (\text{space } M) \{\text{indicator } b - 'A \cap \text{space } M \mid A. A \in \text{borel}\} \subseteq G$ 
    using borel-measurable-indicator[OF A(3), THEN measurable-sets] sets.top
    subalgebra
    by (intro sets.sigma-sets-subset') (fastforce simp add: subalgebra-def)+

    have Un-in-sigma:  $F \cup \text{vimage-algebra } (\text{space } M) \ f \ \text{borel} \subseteq \text{sigma } (\text{space } M) \ (F \cup \text{vimage-algebra } (\text{space } M) \ f \ \text{borel})$  by (metis equalityE le-supI sets.space-closed
    sigma-le-sets space-vimage-algebra subalg subalgebra-def)

    have [intro]: indep-var borel (indicator b) borel  $(\lambda \omega. \text{indicator } a \ \omega *_{\mathbb{R}} f \ \omega)$ 
    proof -
      have [simp]:  $\text{sigma-sets } (\text{space } M) \{(\lambda \omega. \text{indicator } a \ \omega *_{\mathbb{R}} f \ \omega) - 'A \cap \text{space } M \mid A. A \in \text{borel}\} \subseteq \text{sigma } (\text{space } M) \ (F \cup \text{vimage-algebra } (\text{space } M) \ f \ \text{borel})$ 
      proof -
        have *:  $(\lambda \omega. \text{indicator } a \ \omega *_{\mathbb{R}} f \ \omega) \in \text{borel-measurable } (\text{sigma } (\text{space } M) \ (F \cup \text{vimage-algebra } (\text{space } M) \ f \ \text{borel}))$ 
        using borel-measurable-indicator[OF A(2), THEN measurable-sets, OF
        borel-open] subalgebra
        by (intro borel-measurable-scaleR borel-measurableI Un-in-sigma[THEN
        subsetD])
        (auto simp add: space-measure-of-conv subalgebra-def sets-vimage-algebra2)
      thus ?thesis using measurable-sets[OF *] by (intro sets.sigma-sets-subset',

```

```

auto simp add: space-measure-of-conv)
qed
have indep-set (sigma-sets (space M) {indicator b - ' A ∩ space M | A. A ∈
borel}) (sigma-sets (space M) {(λω. indicator a ω *R f ω) - ' A ∩ space M | A. A
∈ borel})
using independent unfolding indep-set-def by (rule indep-sets-mono-sets,
auto split: bool.split)
thus ?thesis by (subst indep-var-eq, auto intro!: borel-measurable-scaleR)
qed

have [intro]: indep-var borel (indicator b) borel (λω. indicat-real a ω *R
cond-exp M F f ω)
proof -
have [simp]: sigma-sets (space M) {(λω. indicator a ω *R cond-exp M F f
ω) - ' A ∩ space M | A. A ∈ borel} ⊆ sigma (space M) (F ∪ vimage-algebra (space
M) f borel)
proof -
have *: (λω. indicator a ω *R cond-exp M F f ω) ∈ borel-measurable (sigma
(space M) (F ∪ vimage-algebra (space M) f borel))
using borel-measurable-indicator[OF A(2), THEN measurable-sets, OF
borel-open] subalgebra
borel-measurable-cond-exp[THEN measurable-sets, OF borel-open, of
- M F f]
by (intro borel-measurable-scaleR borel-measurableI Un-in-sigma[THEN
subsetD])
(auto simp add: space-measure-of-conv subalgebra-def)
thus ?thesis using measurable-sets[OF *] by (intro sets.sigma-sets-subset',
auto simp add: space-measure-of-conv)
qed
have indep-set (sigma-sets (space M) {indicator b - ' A ∩ space M | A. A ∈
borel}) (sigma-sets (space M) {(λω. indicator a ω *R cond-exp M F f ω) - ' A ∩
space M | A. A ∈ borel})
using independent unfolding indep-set-def by (rule indep-sets-mono-sets,
auto split: bool.split)
thus ?thesis by (subst indep-var-eq, auto intro!: borel-measurable-scaleR)
qed

have set-lebesgue-integral M A f = (LINT x|M. indicator b x * (indicator a
x *R f x))
unfolding set-lebesgue-integral-def A indicator-inter-arith
by (intro Bochner-Integration.integral-cong, auto simp add: scaleR-scaleR[symmetric]
indicator-times-eq-if(1))
also have ... = (LINT x|M. indicator b x) * (LINT x|M. indicator a x *R f
x)
by (intro indep-var-lebesgue-integral
Bochner-Integration.integrable-bound[OF integrable-const[of 1 :: 'b]
borel-measurable-indicator]
integrable-mult-indicator[OF - assms(4)], blast) (auto simp add:
indicator-def)

```

```

    also have ... = (LINT x|M. indicator b x) * (LINT x|M. indicator a x *R
cond-exp M F f x)
    using cond-exp-set-integral[OF assms(4) A(2)] unfolding set-lebesgue-integral-def
by argo
    also have ... = (LINT x|M. indicator b x * (indicator a x *R cond-exp M
F f x))
    by (intro indep-var-lebesgue-integral[symmetric]
        Bochner-Integration.integrable-bound[OF integrable-const[of 1 :: 'b]
borel-measurable-indicator]
        integrable-mult-indicator[OF - integrable-cond-exp], blast) (auto simp
add: indicator-def)
    also have ... = set-lebesgue-integral M A (cond-exp M F f)
    unfolding set-lebesgue-integral-def A indicator-inter-arith
    by (intro Bochner-Integration.integrable-cong, auto simp add: scaleR-scaleR[symmetric]
indicator-times-eq-if(1))
    finally show ?case .
next
    case empty
    then show ?case unfolding set-lebesgue-integral-def by simp
next
    case (compl A)
    have A-in-space: A ⊆ space M using compl using in-events sets.sets-into-space
by blast
    have set-lebesgue-integral M (space M - A) f = set-lebesgue-integral M (space
M - A ∪ A) f - set-lebesgue-integral M A f
    using compl(1) in-events
    by (subst set-integral-Un[of space M - A A], blast)
        (simp | intro integrable-mult-indicator[folded set-integrable-def, OF -
assms(4)], fast)+
    also have ... = set-lebesgue-integral M (space M - A ∪ A) (cond-exp M F f)
- set-lebesgue-integral M A (cond-exp M F f)
    using cond-exp-set-integral[OF assms(4) sets.top] compl subalgebra by (simp
add: subalgebra-def Un-absorb2[OF A-in-space])
    also have ... = set-lebesgue-integral M (space M - A) (cond-exp M F f)
    using compl(1) in-events
    by (subst set-integral-Un[of space M - A A], blast)
        (simp | intro integrable-mult-indicator[folded set-integrable-def, OF -
integrable-cond-exp], fast)+
    finally show ?case .
next
    case (union A)
    have set-lebesgue-integral M (⋃ (range A)) f = (∑ i. set-lebesgue-integral M
(A i) f)
    using union in-events
    by (intro lebesgue-integral-countable-add) (auto simp add: disjoint-family-onD
intro!: integrable-mult-indicator[folded set-integrable-def, OF - assms(4)])
    also have ... = (∑ i. set-lebesgue-integral M (A i) (cond-exp M F f)) using
union by presburger
    also have ... = set-lebesgue-integral M (⋃ (range A)) (cond-exp M F f)

```

```

    using union in-events
    by (intro lebesgue-integral-countable-add[symmetric]) (auto simp add: disjoint-family-onD intro!: integrable-mult-indicator[folded set-integrable-def, OF - integrable-cond-exp])
    finally show ?case .
  qed
}
moreover have sigma (space M) {a ∩ b | a b. a ∈ F ∧ b ∈ G} = sigma (space M) (F ∪ G)
proof -
  have sigma-sets (space M) {a ∩ b | a b. a ∈ sets F ∧ b ∈ sets G} = sigma-sets (space M) (sets F ∪ sets G)
  proof -
    {
      fix a b assume asm: a ∈ F b ∈ G
      hence a ∩ b ∈ sigma-sets (space M) (F ∪ G) using subalgebra unfolding Int-range-binary by (intro sigma-sets-Inter[OF - binary-in-sigma-sets]) (force simp add: subalgebra-def dest: sets.sets-into-space)+
    }
    moreover
    {
      fix a
      assume a ∈ sets F
      hence a ∈ sigma-sets (space M) {a ∩ b | a b. a ∈ sets F ∧ b ∈ sets G}
      using subalgebra sets.top[of G] sets.sets-into-space[of - F]
      by (intro sigma-sets.Basic, auto simp add: subalgebra-def)
    }
    moreover
    {
      fix a assume a ∈ sets F ∨ a ∈ sets G a ∉ sets F
      hence a ∈ sets G by blast
      hence a ∈ sigma-sets (space M) {a ∩ b | a b. a ∈ sets F ∧ b ∈ sets G}
      using subalgebra sets.top[of F] sets.sets-into-space[of - G]
      by (intro sigma-sets.Basic, auto simp add: subalgebra-def)
    }
  }
  ultimately show ?thesis by (intro sigma-sets-eqI) auto
qed
thus ?thesis using subalgebra by (intro sigma-eqI) (force simp add: subalgebra-def dest: sets.sets-into-space)+
qed
moreover have (cond-exp M F f) ∈ borel-measurable (sigma (space M) (sets F ∪ sets G))
proof -
  have F ⊆ sigma (space M) (F ∪ G) by (metis Un-least Un-upper1 measure-of-of-measure sets.space-closed sets-measure-of sigma-sets-subseteq subalg subalgebra(2) subalgebra-def)
  thus ?thesis using borel-measurable-cond-exp[THEN measurable-sets, OF borel-open, of - M F f] subalgebra by (intro borel-measurableI, force simp only: space-measure-of-conv subalgebra-def)

```



```

qed
ultimately show ?thesis using assms(4) integrable-cond-exp by (intro Un-sigma.cond-exp-charact)
presburger+
qed

```

If a random variable is independent of a σ -algebra F , its conditional expectation $\text{cond-exp } M F f$ is just its expectation.

```

lemma (in prob-space) cond-exp-indep:
  fixes f :: 'a  $\Rightarrow$  'b :: {second-countable-topology, banach, real-normed-field}
  assumes subalgebra: subalgebra M F
    and independent: indep-set F (vimage-algebra (space M) f borel)
    and integrable: integrable M f
  shows AE x in M. cond-exp M F f x = expectation f
proof -
  have indep-set F (sigma (space M) (sigma (space M) {})  $\cup$  (vimage-algebra (space M) f borel)))
    using independent unfolding indep-set-def
  by (rule indep-sets-mono-sets, simp add: bool.split)
  (metis bot.extremum dual-order.refl sets.sets-measure-of-eq sets.sigma-sets-subset'
  sets-vimage-algebra-space space-vimage-algebra sup.absorb-iff2)
  hence cond-exp-indep: AE x in M. cond-exp M (sigma (space M) (sigma (space M) {}  $\cup$  F)) f x = expectation f
    using cond-exp-indep-subalgebra[OF - subalgebra - integrable, of sigma (space M) {}] cond-exp-trivial[OF integrable]
  by (auto simp add: subalgebra-def sigma-sets-empty-eq)
  have sets (sigma (space M) (sigma (space M) {}  $\cup$  F)) = F
    using subalgebra sets.top[of F] unfolding subalgebra-def
  by (simp add: sigma-sets-empty-eq, subst insert-absorb[of space M F], blast)
  (metis insert-absorb[OF sets.empty-sets] sets.sets-measure-of-eq)
  hence AE x in M. cond-exp M (sigma (space M) (sigma (space M) {}  $\cup$  F)) f
    x = cond-exp M F f x by (rule cond-exp-sets-cong)
  thus ?thesis using cond-exp-indep by force
qed
end

```

```

theory Filtered-Measure
  imports HOL-Probability.Conditional-Expectation
begin

```

5 Filtered Measure Spaces

5.1 Filtered Measure

```

locale filtered-measure =
  fixes M F and t0 :: 'b :: {second-countable-topology, order-topology, t2-space}
  assumes subalgebras:  $\bigwedge i. t_0 \leq i \implies \text{subalgebra } M (F i)$ 
    and sets-F-mono:  $\bigwedge i j. t_0 \leq i \implies i \leq j \implies \text{sets } (F i) \leq \text{sets } (F j)$ 

```

begin

lemma *space-F[simp]*:
assumes $t_0 \leq i$
shows $\text{space } (F i) = \text{space } M$
using *subalgebras assms* **by** (*simp add: subalgebra-def*)

lemma *subalgebra-F[intro]*:
assumes $t_0 \leq i \leq j$
shows $\text{subalgebra } (F j) (F i)$
unfolding *subalgebra-def* **using** *assms* **by** (*simp add: sets-F-mono*)

lemma *borel-measurable-mono*:
assumes $t_0 \leq i \leq j$
shows $\text{borel-measurable } (F i) \subseteq \text{borel-measurable } (F j)$
unfolding *subset-iff* **by** (*metis assms subalgebra-F measurable-from-subalg*)

end

locale *linearly-filtered-measure* = *filtered-measure* $M F t_0$ **for** M **and** $F :: - \Rightarrow \{ \text{linorder-topology} \}$ $\Rightarrow -$ **and** t_0

locale *nat-filtered-measure* = *linearly-filtered-measure* $M F 0$ **for** M **and** $F :: \text{nat} \Rightarrow -$

locale *real-filtered-measure* = *linearly-filtered-measure* $M F 0$ **for** M **and** $F :: \text{real} \Rightarrow -$

5.2 σ -Finite Filtered Measure

The locale presented here is a generalization of the *sigma-finite-subalgebra* for a particular filtration.

locale *sigma-finite-filtered-measure* = *filtered-measure* +
assumes *sigma-finite-initial*: *sigma-finite-subalgebra* $M (F t_0)$

lemma (**in** *sigma-finite-filtered-measure*) *sigma-finite-subalgebra-F[intro]*:
assumes $t_0 \leq i$
shows *sigma-finite-subalgebra* $M (F i)$
using *assms* **by** (*metis dual-order.refl sets-F-mono sigma-finite-initial sigma-finite-subalgebra.nested-subalg-is-subalgebras subalgebra-def*)

locale *nat-sigma-finite-filtered-measure* = *sigma-finite-filtered-measure* $M F 0 :: \text{nat}$ **for** $M F$

locale *real-sigma-finite-filtered-measure* = *sigma-finite-filtered-measure* $M F 0 :: \text{real}$ **for** $M F$

sublocale *nat-sigma-finite-filtered-measure* \subseteq *sigma-finite-subalgebra* $M F i$ **by** *blast*

sublocale *real-sigma-finite-filtered-measure* \subseteq *sigma-finite-subalgebra* $M F |i|$ **by** *fastforce*

5.3 Finite Filtered Measure

locale *finite-filtered-measure* = *filtered-measure* + *finite-measure*

sublocale *finite-filtered-measure* \subseteq *sigma-finite-filtered-measure*

using *subalgebras* **by** (*unfold-locale*, *blast*, *meson* *dual-order.refl* *finite-measure-axioms* *finite-measure-def* *finite-measure-restr-to-subalg* *sigma-finite-measure.sigma-finite-countable*)

locale *nat-finite-filtered-measure* = *finite-filtered-measure* *M F* 0 :: *nat* **for** *M F*

locale *real-finite-filtered-measure* = *finite-filtered-measure* *M F* 0 :: *real* **for** *M F*

sublocale *nat-finite-filtered-measure* \subseteq *nat-sigma-finite-filtered-measure* ..

sublocale *real-finite-filtered-measure* \subseteq *real-sigma-finite-filtered-measure* ..

5.4 Constant Filtration

lemma *filtered-measure-constant-filtration*:

assumes *subalgebra* *M F*

shows *filtered-measure* *M* ($\lambda\cdot. F$) *t*₀

using *assms* **by** (*unfold-locale*) *blast*+

sublocale *sigma-finite-subalgebra* \subseteq *constant-filtration: sigma-finite-filtered-measure*

M $\lambda\cdot. F$:: '*t* :: {*second-countable-topology*, *linorder-topology*}. *F* *t*₀

using *subalg* **by** (*unfold-locale*) *blast*+

lemma (*in* *finite-measure*) *filtered-measure-constant-filtration*:

assumes *subalgebra* *M F*

shows *finite-filtered-measure* *M* ($\lambda\cdot. F$) *t*₀

using *assms* **by** (*unfold-locale*) *blast*+

end

theory *Stochastic-Process*

imports *Filtered-Measure* *Measure-Space-Supplement* *HOL-Probability.Independent-Family*

begin

6 Stochastic Processes

6.1 Stochastic Process

A stochastic process is a collection of random variables, indexed by a type '*b*.

locale *stochastic-process* =

fixes *M* *t*₀ **and** *X* :: '*b* :: {*second-countable-topology*, *order-topology*, *t2-space*} \Rightarrow '*a* \Rightarrow '*c* :: {*second-countable-topology*, *banach*}

assumes *random-variable[measurable]*: $\bigwedge i. t_0 \leq i \implies X\ i \in \text{borel-measurable } M$

begin

definition *left-continuous* **where** *left-continuous* = (*AE* ξ in M . $\forall t$. *continuous* (*at-left* t) (λi . $X\ i\ \xi$))

definition *right-continuous* **where** *right-continuous* = (*AE* ξ in M . $\forall t$. *continuous* (*at-right* t) (λi . $X\ i\ \xi$))

end

We specify the following locales to formalize discrete time and continuous time processes.

locale *nat-stochastic-process* = *stochastic-process* $M\ 0 :: \text{nat } X$ **for** $M\ X$

locale *real-stochastic-process* = *stochastic-process* $M\ 0 :: \text{real } X$ **for** $M\ X$

lemma *stochastic-process-const-fun*:

assumes $f \in \text{borel-measurable } M$

shows *stochastic-process* $M\ t_0\ (\lambda -. f)$ **using** *assms* **by** (*unfold-locales*)

lemma *stochastic-process-const*:

shows *stochastic-process* $M\ t_0\ (\lambda i -. c\ i)$ **by** (*unfold-locales*) *simp*

In the following segment, we cover basic operations on stochastic processes.

context *stochastic-process*

begin

lemma *compose-stochastic*:

assumes $\bigwedge i. t_0 \leq i \implies f\ i \in \text{borel-measurable borel}$

shows *stochastic-process* $M\ t_0\ (\lambda i\ \xi. (f\ i)\ (X\ i\ \xi))$

by (*unfold-locales*) (*intro measurable-compose[OF random-variable assms]*)

lemma *norm-stochastic*: *stochastic-process* $M\ t_0\ (\lambda i\ \xi. \text{norm } (X\ i\ \xi))$ **by** (*fastforce intro: compose-stochastic*)

lemma *scaleR-right-stochastic*:

assumes *stochastic-process* $M\ t_0\ Y$

shows *stochastic-process* $M\ t_0\ (\lambda i\ \xi. (Y\ i\ \xi) *_{\mathbb{R}} (X\ i\ \xi))$

using *stochastic-process.random-variable[OF assms]* *random-variable* **by** (*unfold-locales*) *simp*

lemma *scaleR-right-const-fun-stochastic*:

assumes $f \in \text{borel-measurable } M$

shows *stochastic-process* $M\ t_0\ (\lambda i\ \xi. f\ \xi *_{\mathbb{R}} (X\ i\ \xi))$

by (*unfold-locales*) (*intro borel-measurable-scaleR assms random-variable*)

lemma *scaleR-right-const-stochastic*: *stochastic-process* $M\ t_0\ (\lambda i\ \xi. c\ i *_{\mathbb{R}} (X\ i\ \xi))$

by (*unfold-locales*) *simp*

lemma *add-stochastic*:

assumes *stochastic-process* $M\ t_0\ Y$

shows *stochastic-process* $M\ t_0\ (\lambda i\ \xi. X\ i\ \xi + Y\ i\ \xi)$

using *stochastic-process.random-variable*[*OF assms*] *random-variable* **by** (*unfold-locales*)
simp

lemma *diff-stochastic*:

assumes *stochastic-process* $M\ t_0\ Y$

shows *stochastic-process* $M\ t_0\ (\lambda i\ \xi.\ X\ i\ \xi - Y\ i\ \xi)$

using *stochastic-process.random-variable*[*OF assms*] *random-variable* **by** (*unfold-locales*)
simp

lemma *uminus-stochastic*: *stochastic-process* $M\ t_0\ (-X)$ **using** *scaleR-right-const-stochastic*[*of*
 $\lambda\cdot. -1$] **by** (*simp add: fun-Compl-def*)

lemma *partial-sum-stochastic*: *stochastic-process* $M\ t_0\ (\lambda n\ \xi.\ \sum_{i \in \{t_0..n\}} X\ i\ \xi)$
by (*unfold-locales*) *simp*

lemma *partial-sum'-stochastic*: *stochastic-process* $M\ t_0\ (\lambda n\ \xi.\ \sum_{i \in \{t_0..<n\}} X\ i\ \xi)$
by (*unfold-locales*) *simp*

end

lemma *stochastic-process-sum*:

assumes $\bigwedge i. i \in I \implies \text{stochastic-process } M\ t_0\ (X\ i)$

shows *stochastic-process* $M\ t_0\ (\lambda k\ \xi.\ \sum_{i \in I} X\ i\ k\ \xi)$ **using** *assms*[*THEN*
stochastic-process.random-variable] **by** (*unfold-locales, auto*)

6.1.1 Natural Filtration

The natural filtration induced by a stochastic process X is the filtration generated by all events involving the process up to the time index t , i.e. $F_t = \sigma(\{X\ s \mid s \leq t\})$.

definition *natural-filtration* :: 'a measure \Rightarrow 'b \Rightarrow ('b \Rightarrow 'a \Rightarrow 'c :: *topological-space*) \Rightarrow 'b :: {*second-countable-topology, order-topology*} \Rightarrow 'a measure **where**
natural-filtration $M\ t_0\ Y = (\lambda t.\ \text{family-vimage-algebra } (\text{space } M)\ \{Y\ i \mid i. i \in \{t_0..t\}\}\ \text{borel})$

abbreviation *nat-natural-filtration* $\equiv \lambda M.\ \text{natural-filtration } M\ (0 :: \text{nat})$

abbreviation *real-natural-filtration* $\equiv \lambda M.\ \text{natural-filtration } M\ (0 :: \text{real})$

lemma *space-natural-filtration*[*simp*]: *space* (*natural-filtration* $M\ t_0\ X\ t$) = *space*
 M **unfolding** *natural-filtration-def* *space-family-vimage-algebra* ..

lemma *sets-natural-filtration*: *sets* (*natural-filtration* $M\ t_0\ X\ t$) = *sigma-sets* (*space*
 M) ($\bigcup_{i \in \{t_0..t\}} \{X\ i - 'A \cap \text{space } M \mid A. A \in \text{borel}\}$)

unfolding *natural-filtration-def* *sets-family-vimage-algebra* **by** (*intro sigma-sets-eqI*)
blast+

lemma *sets-natural-filtration'*:

assumes *borel* = *sigma UNIV* S

shows $\text{sets } (\text{natural-filtration } M \ t_0 \ X \ t) = \text{sigma-sets } (\text{space } M) \ (\bigcup_{i \in \{t_0..t\}}. \{X \ i - ' A \cap \text{space } M \mid A. A \in S\})$
proof (*subst sets-natural-filtration, intro sigma-sets-eqI, clarify*)
fix i **and** $A :: 'a \text{ set}$ **assume** $\text{asm: } i \in \{t_0..t\} \ A \in \text{sets borel}$
hence $A \in \text{sigma-sets UNIV } S$ **unfolding** *assms* **by** *simp*
thus $X \ i - ' A \cap \text{space } M \in \text{sigma-sets } (\text{space } M) \ (\bigcup_{i \in \{t_0..t\}}. \{X \ i - ' A \cap \text{space } M \mid A. A \in S\})$
proof (*induction*)
case (*Compl a*)
have $X \ i - ' (UNIV - a) \cap \text{space } M = \text{space } M - (X \ i - ' a \cap \text{space } M)$ **by** *blast*
then show *?case* **using** *Compl(2)[THEN sigma-sets.Compl]* **by** *presburger*
next
case (*Union a*)
have $X \ i - ' \bigcup (\text{range } a) \cap \text{space } M = \bigcup (\text{range } (\lambda j. X \ i - ' a \ j \cap \text{space } M))$
by *blast*
then show *?case* **using** *Union(2)[THEN sigma-sets.Union]* **by** *presburger*
qed (*auto intro: asm sigma-sets.Empty*)
qed (*intro sigma-sets.Basic, force simp add: assms*)

lemma *sets-natural-filtration-open:*

$\text{sets } (\text{natural-filtration } M \ t_0 \ X \ t) = \text{sigma-sets } (\text{space } M) \ (\bigcup_{i \in \{t_0..t\}}. \{X \ i - ' A \cap \text{space } M \mid A. \text{open } A\})$
using *sets-natural-filtration'* **by** (*force simp only: borel-def mem-Collect-eq*)

lemma *sets-natural-filtration-oi:*

$\text{sets } (\text{natural-filtration } M \ t_0 \ X \ t) = \text{sigma-sets } (\text{space } M) \ (\bigcup_{i \in \{t_0..t\}}. \{X \ i - ' A \cap \text{space } M \mid A :: - :: \{\text{linorder-topology, second-countable-topology}\} \text{ set. } A \in \text{range greaterThan}\})$
by (*rule sets-natural-filtration'[OF borel-Ioi]*)

lemma *sets-natural-filtration-io:*

$\text{sets } (\text{natural-filtration } M \ t_0 \ X \ t) = \text{sigma-sets } (\text{space } M) \ (\bigcup_{i \in \{t_0..t\}}. \{X \ i - ' A \cap \text{space } M \mid A :: - :: \{\text{linorder-topology, second-countable-topology}\} \text{ set. } A \in \text{range lessThan}\})$
by (*rule sets-natural-filtration'[OF borel-Iio]*)

lemma *sets-natural-filtration-ci:*

$\text{sets } (\text{natural-filtration } M \ t_0 \ X \ t) = \text{sigma-sets } (\text{space } M) \ (\bigcup_{i \in \{t_0..t\}}. \{X \ i - ' A \cap \text{space } M \mid A :: \text{real set. } A \in \text{range atLeast}\})$
by (*rule sets-natural-filtration'[OF borel-Ici]*)

context *stochastic-process*

begin

lemma *subalgebra-natural-filtration:*

shows *subalgebra* $M \ (\text{natural-filtration } M \ t_0 \ X \ i)$
unfolding *subalgebra-def* **using** *measurable-family-iff-sets* **by** (*force simp add: natural-filtration-def*)

lemma *filtered-measure-natural-filtration:*

shows *filtered-measure* M (*natural-filtration* $M \ t_0 \ X$) t_0
by (*unfold-locales*) (*intro subalgebra-natural-filtration, simp only: sets-natural-filtration, intro sigma-sets-subseteq, force*)

In order to show that the natural filtration constitutes a filtered σ -finite measure, we need to provide a countable exhausting set in the preimage of $X \ t_0$.

lemma *sigma-finite-filtered-measure-natural-filtration:*

assumes *exhausting-set: countable* $A \ (\bigcup A) = \text{space } M \ \bigwedge a. a \in A \implies \text{emeasure } M \ a \neq \infty \ \bigwedge a. a \in A \implies \exists b \in \text{borel}. a = X \ t_0 - ' b \cap \text{space } M$

shows *sigma-finite-filtered-measure* M (*natural-filtration* $M \ t_0 \ X$) t_0

proof (*unfold-locales*)

have $A \subseteq \text{sets } (\text{restr-to-subalg } M \ (\text{natural-filtration } M \ t_0 \ X \ t_0))$ **using** *exhausting-set* **by** (*simp add: sets-restr-to-subalg[OF subalgebra-natural-filtration] sets-natural-filtration*) *fast*

moreover have $\bigcup A = \text{space } (\text{restr-to-subalg } M \ (\text{natural-filtration } M \ t_0 \ X \ t_0))$

unfolding *space-restr-to-subalg* **using** *exhausting-set* **by** *simp*

moreover have $\forall a \in A. \text{emeasure } (\text{restr-to-subalg } M \ (\text{natural-filtration } M \ t_0 \ X \ t_0)) \ a \neq \infty$ **using** *calculation(1) exhausting-set(3)*

by (*auto simp add: sets-restr-to-subalg[OF subalgebra-natural-filtration] emeasure-restr-to-subalg[OF subalgebra-natural-filtration]*)

ultimately show $\exists A. \text{countable } A \wedge A \subseteq \text{sets } (\text{restr-to-subalg } M \ (\text{natural-filtration } M \ t_0 \ X \ t_0)) \wedge \bigcup A = \text{space } (\text{restr-to-subalg } M \ (\text{natural-filtration } M \ t_0 \ X \ t_0)) \wedge (\forall a \in A. \text{emeasure } (\text{restr-to-subalg } M \ (\text{natural-filtration } M \ t_0 \ X \ t_0)) \ a \neq \infty)$ **using** *exhausting-set* **by** *blast*

show $\bigwedge i j. \llbracket t_0 \leq i; i \leq j \rrbracket \implies \text{sets } (\text{natural-filtration } M \ t_0 \ X \ i) \subseteq \text{sets } (\text{natural-filtration } M \ t_0 \ X \ j)$ **using** *filtered-measure.subalgebra-F[OF filtered-measure-natural-filtration]*

by (*simp add: subalgebra-def*)

qed (*auto intro: subalgebra-natural-filtration*)

lemma *finite-filtered-measure-natural-filtration:*

assumes *finite-measure* M

shows *finite-filtered-measure* M (*natural-filtration* $M \ t_0 \ X$) t_0

using *finite-measure.axioms[OF assms] filtered-measure-natural-filtration* **by** *intro-locales*

end

Filtration generated by independent variables.

lemma (*in prob-space*) *indep-set-natural-filtration:*

assumes $t_0 \leq s \leq t$ *indep-vars* $(\lambda-. \text{borel}) \ X \ \{t_0..s\}$

shows *indep-set* (*natural-filtration* $M \ t_0 \ X \ s$) (*vimage-algebra* (*space* M) ($X \ t$) *borel*)

proof –

have *indep-sets* $(\lambda i. \{X \ i - ' A \cap \text{space } M \mid A. A \in \text{sets borel}\}) \ (\bigcup (\text{range } (\text{case-bool } \{t_0..s\} \ \{t\})))$

using *assms*

```

by (intro assms( $\beta$ )[unfolded indep-vars-def, THEN conjunct2, THEN indep-sets-mono])
(auto simp add: case-bool-if)
thus ?thesis unfolding indep-set-def using assms
by (intro indep-sets-cong[THEN iffD1, OF refl - indep-sets-collect-sigma[of  $\lambda i.$ 
 $\{X\ i - 'A \cap \text{space } M \mid A. A \in \text{borel}\}$  case-bool  $\{t_0..s\}$   $\{t\}$ ]])
(simp add: sets-natural-filtration sets-vimage-algebra split: bool.split, simp,
intro Int-stableI, clarsimp, metis sets.Int vimage-Int Int-commute Int-left-absorb
Int-left-commute, force simp add: disjoint-family-on-def split: bool.split)
qed

```

6.2 Adapted Process

We call a collection a stochastic process X adapted if $X\ i$ is $F\ i$ -borel-measurable for all indices i .

```

locale adapted-process = filtered-measure  $M\ F\ t_0$  for  $M\ F\ t_0$  and  $X :: - \Rightarrow - \Rightarrow -$ 
:: {second-countable-topology, banach} +
assumes adapted[measurable]:  $\bigwedge i. t_0 \leq i \implies X\ i \in \text{borel-measurable } (F\ i)$ 
begin

```

```

lemma adaptedE[elim]:
assumes  $\llbracket \bigwedge j\ i. t_0 \leq j \implies j \leq i \implies X\ j \in \text{borel-measurable } (F\ i) \rrbracket \implies P$ 
shows  $P$ 
using assms using adapted by (metis dual-order.trans borel-measurable-subalgebra
sets-F-mono space-F)

```

```

lemma adaptedD:
assumes  $t_0 \leq j\ j \leq i$ 
shows  $X\ j \in \text{borel-measurable } (F\ i)$  using assms adaptedE by meson

```

end

```

locale nat-adapted-process = adapted-process  $M\ F\ 0 :: \text{nat } X$  for  $M\ F\ X$ 
locale real-adapted-process = adapted-process  $M\ F\ 0 :: \text{real } X$  for  $M\ F\ X$ 

```

```

sublocale nat-adapted-process  $\subseteq$  nat-filtered-measure ..
sublocale real-adapted-process  $\subseteq$  real-filtered-measure ..

```

```

lemma (in filtered-measure) adapted-process-const-fun:
assumes  $f \in \text{borel-measurable } (F\ t_0)$ 
shows adapted-process  $M\ F\ t_0\ (\lambda -. f)$ 
using measurable-from-subalg subalgebra-F assms by (unfold-locales) blast

```

```

lemma (in filtered-measure) adapted-process-const:
shows adapted-process  $M\ F\ t_0\ (\lambda i -. c\ i)$  by (unfold-locales) simp

```

Again, we cover basic operations.

```

context adapted-process
begin

```


lemma *compose-adapted*:

assumes $\bigwedge i. t_0 \leq i \implies f\ i \in \text{borel-measurable borel}$
shows *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. (f\ i)\ (X\ i\ \xi))$
by (*unfold-locales*) (*intro measurable-compose*[*OF adapted assms*])

lemma *norm-adapted*: *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. \text{norm}\ (X\ i\ \xi))$ **by** (*fastforce intro: compose-adapted*)

lemma *scaleR-right-adapted*:

assumes *adapted-process* $M\ F\ t_0\ R$
shows *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. (R\ i\ \xi) *_{\mathbb{R}} (X\ i\ \xi))$
using *adapted-process.adapted*[*OF assms*] **adapted by** (*unfold-locales*) *simp*

lemma *scaleR-right-const-fun-adapted*:

assumes $f \in \text{borel-measurable}\ (F\ t_0)$
shows *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. f\ \xi *_{\mathbb{R}} (X\ i\ \xi))$
using *assms* **by** (*fast intro: scaleR-right-adapted adapted-process-const-fun*)

lemma *scaleR-right-const-adapted*: *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. c\ i *_{\mathbb{R}} (X\ i\ \xi))$
by (*unfold-locales*) *simp*

lemma *add-adapted*:

assumes *adapted-process* $M\ F\ t_0\ Y$
shows *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi + Y\ i\ \xi)$
using *adapted-process.adapted*[*OF assms*] **adapted by** (*unfold-locales*) *simp*

lemma *diff-adapted*:

assumes *adapted-process* $M\ F\ t_0\ Y$
shows *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi - Y\ i\ \xi)$
using *adapted-process.adapted*[*OF assms*] **adapted by** (*unfold-locales*) *simp*

lemma *uminus-adapted*: *adapted-process* $M\ F\ t_0\ (-X)$ **using** *scaleR-right-const-adapted*[*of $\lambda -. -1$*] **by** (*simp add: fun-Compl-def*)

lemma *partial-sum-adapted*: *adapted-process* $M\ F\ t_0\ (\lambda n\ \xi. \sum_{i \in \{t_0..n\}} X\ i\ \xi)$

proof (*unfold-locales*)

fix $i :: 'b$
have $X\ j \in \text{borel-measurable}\ (F\ i)$ **if** $t_0 \leq j$ **for** j **using** *that adaptedE* **by** *meson*
thus $(\lambda \xi. \sum_{i \in \{t_0..i\}} X\ i\ \xi) \in \text{borel-measurable}\ (F\ i)$ **by** *simp*
qed

lemma *partial-sum'-adapted*: *adapted-process* $M\ F\ t_0\ (\lambda n\ \xi. \sum_{i \in \{t_0..<n\}} X\ i\ \xi)$

proof (*unfold-locales*)

fix $i :: 'b$
have $X\ j \in \text{borel-measurable}\ (F\ i)$ **if** $t_0 \leq j < i$ **for** j **using** *that adaptedE* **by** *fastforce*

thus $(\lambda \xi. \sum_{i \in \{t_0..<i\}} X \ i \ \xi) \in \text{borel-measurable } (F \ i)$ **by** *simp*
qed

end

In the discrete time case, we have the following lemma which will be useful later on.

lemma (in *nat-adapted-process*) *partial-sum-Suc-adapted*: *nat-adapted-process* M
 $F \ (\lambda n \ \xi. \sum_{i < n} X \ (Suc \ i) \ \xi)$

proof (*unfold-locales*)

fix i

have $X \ j \in \text{borel-measurable } (F \ i)$ **if** $j \leq i$ **for** j **using** *that adaptedD* **by** *blast*

thus $(\lambda \xi. \sum_{i < i} X \ (Suc \ i) \ \xi) \in \text{borel-measurable } (F \ i)$ **by** *auto*

qed

lemma (in *filtered-measure*) *adapted-process-sum*:

assumes $\bigwedge i. i \in I \implies \text{adapted-process } M \ F \ t_0 \ (X \ i)$

shows *adapted-process* $M \ F \ t_0 \ (\lambda k \ \xi. \sum_{i \in I} X \ i \ k \ \xi)$

proof –

{

fix $i \ k$ **assume** $i \in I$ **and** *asm*: $t_0 \leq k$

then interpret *adapted-process* $M \ F \ t_0 \ X \ i$ **using** *assms* **by** *simp*

have $X \ i \ k \in \text{borel-measurable } M \ X \ i \ k \in \text{borel-measurable } (F \ k)$ **using** *measurable-from-subalg subalgebras adapted asm* **by** (*blast*, *simp*)

}

thus *?thesis* **by** (*unfold-locales*) *simp*

qed

An adapted process is necessarily a stochastic process.

sublocale *adapted-process* \subseteq *stochastic-process* **using** *measurable-from-subalg subalgebras adapted* **by** (*unfold-locales*) *blast*

sublocale *nat-adapted-process* \subseteq *nat-stochastic-process* ..

sublocale *real-adapted-process* \subseteq *real-stochastic-process* ..

A stochastic process is always adapted to the natural filtration it generates.

lemma (in *stochastic-process*) *adapted-process-natural-filtration*: *adapted-process*
 $M \ (\text{natural-filtration } M \ t_0 \ X) \ t_0 \ X$

using *filtered-measure-natural-filtration*

by (*intro-locales*) (*auto simp add: natural-filtration-def intro!: adapted-process-axioms.intro measurable-family-vimage-algebra*)

6.3 Progressively Measurable Process

locale *progressive-process* = *filtered-measure* $M \ F \ t_0$ **for** $M \ F \ t_0$ **and** $X :: - \Rightarrow -$
 $\Rightarrow - :: \{\text{second-countable-topology, banach}\} +$

assumes *progressive[measurable]*: $\bigwedge t. t_0 \leq t \implies (\lambda(i, x). X \ i \ x) \in \text{borel-measurable}$
(restrict-space borel \{t_0..t\} \otimes_M F t)

begin

lemma *progressiveD*:
assumes $S \in \text{borel}$
shows $(\lambda(j, \xi). X j \xi) - ' S \cap (\{t_0..i\} \times \text{space } M) \in (\text{restrict-space borel } \{t_0..i\} \otimes_M F i)$
using *measurable-sets*[*OF progressive, OF - assms, of i*]
by (*cases* $t_0 \leq i$) (*auto simp add: space-restrict-space sets-pair-measure space-pair-measure*)
end

locale *nat-progressive-process* = *progressive-process* $M F 0 :: \text{nat } X$ **for** $M F X$
locale *real-progressive-process* = *progressive-process* $M F 0 :: \text{real } X$ **for** $M F X$

lemma (*in filtered-measure*) *progressive-process-const-fun*:
assumes $f \in \text{borel-measurable } (F t_0)$
shows *progressive-process* $M F t_0 (\lambda-. f)$
proof (*unfold-locales*)
fix i **assume** *asm*: $t_0 \leq i$
have $f \in \text{borel-measurable } (F i)$ **using** *borel-measurable-mono*[*OF order.refl asm*]
assms **by** *blast*
thus *case-prod* $(\lambda-. f) \in \text{borel-measurable } (\text{restrict-space borel } \{t_0..i\} \otimes_M F i)$
using *measurable-compose*[*OF measurable-snd*] **by** *simp*
qed

lemma (*in filtered-measure*) *progressive-process-const*:
assumes $c \in \text{borel-measurable borel}$
shows *progressive-process* $M F t_0 (\lambda i -. c i)$
using *assms* **by** (*unfold-locales*) (*auto simp add: measurable-split-conv intro!: measurable-compose*[*OF measurable-fst*] *measurable-restrict-space1*)

context *progressive-process*
begin

lemma *compose-progressive*:
assumes *case-prod* $f \in \text{borel-measurable borel}$
shows *progressive-process* $M F t_0 (\lambda i \xi. (f i) (X i \xi))$
proof
fix i **assume** *asm*: $t_0 \leq i$
have $(\lambda(j, \xi). (j, X j \xi)) \in (\text{restrict-space borel } \{t_0..i\} \otimes_M F i) \rightarrow_M \text{borel } \otimes_M \text{borel}$
using *progressive*[*OF asm*] *measurable-fst''*[*OF measurable-restrict-space1, OF measurable-id*]
by (*auto simp add: measurable-pair-iff measurable-split-conv*)
moreover **have** $(\lambda(j, \xi). f j (X j \xi)) = \text{case-prod } f \circ ((\lambda(j, y). (j, y)) \circ (\lambda(j, \xi). (j, X j \xi)))$ **by** *fastforce*
ultimately show $(\lambda(j, \xi). (f j) (X j \xi)) \in \text{borel-measurable } (\text{restrict-space borel } \{t_0..i\} \otimes_M F i)$ **using** *assms* **by** (*simp add: borel-prod*)
qed

lemma *norm-progressive*: *progressive-process* $M F t_0 (\lambda i \xi. \text{norm } (X i \xi))$ **using** *measurable-compose*[*OF progressive borel-measurable-norm*] **by** (*unfold-locales*) *simp*

lemma *scaleR-right-progressive*:
assumes *progressive-process* $M F t_0 R$
shows *progressive-process* $M F t_0 (\lambda i \xi. (R i \xi) *_{\mathbb{R}} (X i \xi))$
using *progressive-process.progressive*[*OF assms*] **by** (*unfold-locales*) (*simp add: progressive assms*)

lemma *scaleR-right-const-fun-progressive*:
assumes $f \in \text{borel-measurable } (F t_0)$
shows *progressive-process* $M F t_0 (\lambda i \xi. f \xi *_{\mathbb{R}} (X i \xi))$
using *assms* **by** (*fast intro: scaleR-right-progressive progressive-process-const-fun*)

lemma *scaleR-right-const-progressive*:
assumes $c \in \text{borel-measurable borel}$
shows *progressive-process* $M F t_0 (\lambda i \xi. c i *_{\mathbb{R}} (X i \xi))$
using *assms* **by** (*fastforce intro: scaleR-right-progressive progressive-process-const*)

lemma *add-progressive*:
assumes *progressive-process* $M F t_0 Y$
shows *progressive-process* $M F t_0 (\lambda i \xi. X i \xi + Y i \xi)$
using *progressive-process.progressive*[*OF assms*] **by** (*unfold-locales*) (*simp add: progressive assms*)

lemma *diff-progressive*:
assumes *progressive-process* $M F t_0 Y$
shows *progressive-process* $M F t_0 (\lambda i \xi. X i \xi - Y i \xi)$
using *progressive-process.progressive*[*OF assms*] **by** (*unfold-locales*) (*simp add: progressive assms*)

lemma *uminus-progressive*: *progressive-process* $M F t_0 (-X)$ **using** *scaleR-right-const-progressive*[*of $\lambda -. -1$*] **by** (*simp add: fun-Compl-def*)

end

A progressively measurable process is also adapted.

sublocale *progressive-process* \subseteq *adapted-process* **using** *measurable-compose-rev*[*OF progressive measurable-Pair1*]
unfolding *prod.case space-restrict-space*
by *unfold-locales simp*

sublocale *nat-progressive-process* \subseteq *nat-adapted-process* ..
sublocale *real-progressive-process* \subseteq *real-adapted-process* ..

In the discrete setting, adaptedness is equivalent to progressive measurability.

theorem *nat-progressive-iff-adapted*: *nat-progressive-process* $M F X \longleftrightarrow$ *nat-adapted-process*

```

M F X
proof (intro iffI)
  assume asm: nat-progressive-process M F X
  interpret nat-progressive-process M F X by (rule asm)
  show nat-adapted-process M F X ..
next
  assume asm: nat-adapted-process M F X
  interpret nat-adapted-process M F X by (rule asm)
  show nat-progressive-process M F X
  proof (unfold-locales, intro borel-measurableI)
    fix S :: 'b set and i :: nat assume open-S: open S
    {
      fix j assume asm: j ≤ i
      hence X j - ' S ∩ space M ∈ F i using adaptedD[of j, THEN measurable-sets]
      space-F open-S by fastforce
      moreover have case-prod X - ' S ∩ {j} × space M = {j} × (X j - ' S ∩
      space M) for j by fast
      moreover have {j :: nat} ∈ restrict-space borel {0..i} using asm by (simp
      add: sets-restrict-space-iff)
      ultimately have case-prod X - ' S ∩ {j} × space M ∈ restrict-space borel
      {0..i} ⊗M F i by simp
    }
    hence (λj. (λ(x, y). X x y) - ' S ∩ {j} × space M) ' {..i} ⊆ restrict-space borel
    {0..i} ⊗M F i by blast
    moreover have case-prod X - ' S ∩ space (restrict-space borel {0..i} ⊗M F
    i) = (⋃ j ≤ i. case-prod X - ' S ∩ {j} × space M) unfolding space-pair-measure
    space-restrict-space space-F by force
    ultimately show case-prod X - ' S ∩ space (restrict-space borel {0..i} ⊗M
    F i) ∈ restrict-space borel {0..i} ⊗M F i by (metis sets.countable-UN)
  qed
qed

```

6.4 Predictable Process

We introduce the constant Σ_P to denote the predictable σ -algebra.

context linearly-filtered-measure
begin

definition $\Sigma_P :: ('b \times 'a) \text{ measure where predictable-sigma: } \Sigma_P \equiv \text{sigma } (\{t_{0..}\} \times \text{space } M) (\{\{s <..t\} \times A \mid A \text{ s t. } A \in F \text{ s } \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in F t_0\})$

lemma space-predictable-sigma[simp]: space $\Sigma_P = (\{t_{0..}\} \times \text{space } M)$ **unfolding** predictable-sigma space-measure-of-conv **by** blast

lemma sets-predictable-sigma: sets $\Sigma_P = \text{sigma-sets } (\{t_{0..}\} \times \text{space } M) (\{\{s <..t\} \times A \mid A \text{ s t. } A \in F \text{ s } \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in F t_0\})$
unfolding predictable-sigma **using** space-F sets.sets-into-space **by** (subst sets-measure-of) fastforce+

lemma *measurable-predictable-sigma-snd*:
assumes *countable* \mathcal{I} $\mathcal{I} \subseteq \{\{s <..t\} \mid s \ t. \ t_0 \leq s \wedge s < t\} \ \{t_0 <..\} \subseteq (\bigcup \mathcal{I})$
shows $snd \in \Sigma_P \rightarrow_M F \ t_0$
proof (*intro measurableI*)
fix $S :: 'a \text{ set}$ **assume** $asm: S \in F \ t_0$
have *countable*: *countable* $((\lambda I. I \times S) \text{ ' } \mathcal{I})$ **using** *assms(1)* **by** *blast*
have $(\lambda I. I \times S) \text{ ' } \mathcal{I} \subseteq \{\{s <..t\} \times A \mid A \ s \ t. A \in F \ s \wedge t_0 \leq s \wedge s < t\}$ **using**
sets-F-mono[*OF order-refl*, *THEN subsetD*, *OF - asm*] *assms(2)* **by** *blast*
hence $(\bigcup I \in \mathcal{I}. I \times S) \cup \{t_0\} \times S \in \Sigma_P$ **unfolding** *sets-predictable-sigma* **using**
asm **by** (*intro sigma-sets-Un*[*OF sigma-sets-UNION*[*OF countable*] *sigma-sets.Basic*]
sigma-sets.Basic) *blast+*
moreover **have** $snd - ' S \cap \text{space } \Sigma_P = \{t_0..\} \times S$ **using** *sets.sets-into-space*[*OF*
asm] **by** *fastforce*
moreover **have** $\{t_0\} \cup \{t_0 <..\} = \{t_0..\}$ **by** *auto*
moreover **have** $(\bigcup I \in \mathcal{I}. I \times S) \cup \{t_0\} \times S = \{t_0..\} \times S$ **using** *assms(2,3)*
calculation(3) **by** *fastforce*
ultimately show $snd - ' S \cap \text{space } \Sigma_P \in \Sigma_P$ **by** *argo*
qed (*auto*)

lemma *measurable-predictable-sigma-fst*:
assumes *countable* \mathcal{I} $\mathcal{I} \subseteq \{\{s <..t\} \mid s \ t. \ t_0 \leq s \wedge s < t\} \ \{t_0 <..\} \subseteq (\bigcup \mathcal{I})$
shows $fst \in \Sigma_P \rightarrow_M \text{borel}$
proof –
have $A \times \text{space } M \in \text{sets } \Sigma_P$ **if** $A \in \text{sigma-sets } \{t_0..\} \ \{\{s <..t\} \mid s \ t. \ t_0 \leq s \wedge s$
 $< t\}$ **for** A **unfolding** *sets-predictable-sigma* **using** *that*
proof (*induction rule: sigma-sets.induct*)
case (*Basic a*)
thus *?case* **using** *space-F sets.top* **by** *blast*
next
case (*Compl a*)
have $(\{t_0..\} - a) \times \text{space } M = \{t_0..\} \times \text{space } M - a \times \text{space } M$ **by** *blast*
then show *?case* **using** *Compl(2)*[*THEN sigma-sets.Compl*] **by** *presburger*
next
case (*Union a*)
have $\bigcup (\text{range } a) \times \text{space } M = \bigcup (\text{range } (\lambda i. a \ i \times \text{space } M))$ **by** *blast*
then show *?case* **using** *Union(2)*[*THEN sigma-sets.Union*] **by** *presburger*
qed (*auto*)
moreover **have** *restrict-space borel* $\{t_0..\} = \text{sigma } \{t_0..\} \ \{\{s <..t\} \mid s \ t. \ t_0 \leq s$
 $\wedge s < t\}$
proof –
have *sigma-sets* $\{t_0..\} \ ((\cap) \ \{t_0..\} \text{ ' } \text{sigma-sets UNIV } (\text{range greaterThan})) =$
sigma-sets $\{t_0..\} \ \{\{s <..t\} \mid s \ t. \ t_0 \leq s \wedge s < t\}$
proof (*intro sigma-sets-eqI* ; *clarify*)
fix $A :: 'b \text{ set}$ **assume** $asm: A \in \text{sigma-sets UNIV } (\text{range greaterThan})$
thus $\{t_0..\} \cap A \in \text{sigma-sets } \{t_0..\} \ \{\{s <..t\} \mid s \ t. \ t_0 \leq s \wedge s < t\}$
proof (*induction rule: sigma-sets.induct*)
case (*Basic a*)
then obtain s **where** $s: a = \{s <..\}$ **by** *blast*

```

show ?case
proof (cases  $t_0 \leq s$ )
  case True
    hence *:  $\{t_0..\} \cap a = (\bigcup i \in \mathcal{I}. \{s<..\} \cap i)$  using  $s$  assms(3) by force
    have  $((\cap) \{s<..\} \text{ ' } \mathcal{I}) \subseteq \text{sigma-sets } \{t_0..\} \{\{s<..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\}$ 
    proof (clarify)
      fix  $A$  assume  $A \in \mathcal{I}$ 
      then obtain  $s' t'$  where  $A: A = \{s'<..t'\}$   $t_0 \leq s' s' < t'$  using assms(2)
    by blast
      hence  $\{s<..\} \cap A = \{\max s s' <..t'\}$  by fastforce
      moreover have  $t_0 \leq \max s s'$  using  $A$  True by linarith
      moreover have  $\max s s' < t'$  if  $s < t'$  using  $A$  that by linarith
      moreover have  $\{s<..\} \cap A = \{\}$  if  $\neg s < t'$  using  $A$  that by force
      ultimately show  $\{s<..\} \cap A \in \text{sigma-sets } \{t_0..\} \{\{s<..t\} \mid s \text{ t. } t_0 \leq s$ 
 $\wedge s < t\}$  by (cases  $s < t'$ ) (blast, simp add: sigma-sets.Empty)
    qed
    thus ?thesis unfolding * using assms(1) by (intro sigma-sets-UNION)
  auto
  next
    case False
    hence  $\{t_0..\} \cap a = \{t_0..\}$  using  $s$  by force
    thus ?thesis using sigma-sets-top by auto
  qed
  next
    case (Compl  $a$ )
    have  $\{t_0..\} \cap (\text{UNIV} - a) = \{t_0..\} - (\{t_0..\} \cap a)$  by blast
    then show ?case using Compl(2)[THEN sigma-sets.Compl] by presburger
  next
    case (Union  $a$ )
    have  $\{t_0..\} \cap \bigcup (\text{range } a) = \bigcup (\text{range } (\lambda i. \{t_0..\} \cap a \ i))$  by blast
    then show ?case using Union(2)[THEN sigma-sets.Union] by presburger
  qed (simp add: sigma-sets.Empty)
  next
    fix  $s \text{ t}$  assume asm:  $t_0 \leq s < t$ 
    hence *:  $\{s<..t\} = \{s<..\} \cap (\{t_0..\} - \{t<..\})$  by force
    have  $\{s<..\} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\} \text{ ' } \text{sigma-sets UNIV (range greaterThan)})$ 
  using asm by (intro sigma-sets.Basic) auto
    moreover have  $\{t_0..\} - \{t<..\} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\} \text{ ' } \text{sigma-sets$ 
 $\text{UNIV (range greaterThan)})$  using asm by (intro sigma-sets.Compl sigma-sets.Basic)
  auto
    ultimately show  $\{s<..t\} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\} \text{ ' } \text{sigma-sets UNIV$ 
 $(\text{range greaterThan}))$  unfolding * Int-range-binary[of  $\{s<..\}$ ] by (intro sigma-sets-Inter[OF
 $- \text{binary-in-sigma-sets}$ ]) auto
  qed
  thus ?thesis unfolding borel-Ioi restrict-space-def emeasure-sigma by (force
intro: sigma-eqI)
  qed
  ultimately have restrict-space borel  $\{t_0..\} \otimes_M \text{sigma (space } M) \{\} \subseteq \text{sets } \Sigma_P$ 

```

```

unfolding sets-pair-measure space-restrict-space space-measure-of-conv
using space-predictable-sigma sets.sigma-algebra-axioms[of  $\Sigma_P$ ]
by (intro sigma-algebra.sigma-sets-subset) (auto simp add: sigma-sets-empty-eq
sets-measure-of-conv)
moreover have space (restrict-space borel  $\{t_0..\}$   $\otimes_M$  sigma (space M)  $\{\}$ ) =
space  $\Sigma_P$  by (simp add: space-pair-measure)
moreover have fst  $\in$  restrict-space borel  $\{t_0..\}$   $\otimes_M$  sigma (space M)  $\{\}$   $\rightarrow_M$ 
borel by (fastforce intro: measurable-fst'[OF measurable-restrict-space1, of  $\lambda x. x$ ])

ultimately show ?thesis by (meson borel-measurable-subalgebra)
qed

end

```

```

locale predictable-process = linearly-filtered-measure M F  $t_0$  for M F  $t_0$  and X ::
-  $\Rightarrow$  -  $\Rightarrow$  - :: {second-countable-topology, banach} +
assumes predictable:  $(\lambda(t, x). X t x) \in$  borel-measurable  $\Sigma_P$ 
begin

```

```

lemmas predictableD = measurable-sets[OF predictable, unfolded space-predictable-sigma]

end

```

```

locale nat-predictable-process = predictable-process M F 0 :: nat X for M F X
locale real-predictable-process = predictable-process M F 0 :: real X for M F X

```

```

lemma (in nat-filtered-measure) measurable-predictable-sigma-snd':
shows snd  $\in \Sigma_P \rightarrow_M F 0$ 
by (intro measurable-predictable-sigma-snd[of range  $(\lambda x. \{Suc x\})$ ]) (force | simp
add: greaterThan-0)+

```

```

lemma (in nat-filtered-measure) measurable-predictable-sigma-fst':
shows fst  $\in \Sigma_P \rightarrow_M$  borel
by (intro measurable-predictable-sigma-fst[of range  $(\lambda x. \{Suc x\})$ ]) (force | simp
add: greaterThan-0)+

```

```

lemma (in real-filtered-measure) measurable-predictable-sigma-snd':
shows snd  $\in \Sigma_P \rightarrow_M F 0$ 
using real-arch-simple by (intro measurable-predictable-sigma-snd[of range  $(\lambda x::nat. \{0 < .. real (Suc x)\})$ ]) (fastforce intro: add-increasing)+

```

```

lemma (in real-filtered-measure) measurable-predictable-sigma-fst':
shows fst  $\in \Sigma_P \rightarrow_M$  borel
using real-arch-simple by (intro measurable-predictable-sigma-fst[of range  $(\lambda x::nat. \{0 < .. real (Suc x)\})$ ]) (fastforce intro: add-increasing)+

```

We show sufficient conditions for functions constant in one argument to

constitute a predictable process. In contrast to the cases before, this is not a triviality.

lemma (in *linearly-filtered-measure*) *predictable-process-const-fun*:
assumes $snd \in \Sigma_P \rightarrow_M F \ t_0 \ f \in \text{borel-measurable } (F \ t_0)$
shows *predictable-process* $M \ F \ t_0 \ (\lambda \cdot. f)$
using *measurable-compose-rev*[*OF* *assms*(2)] *assms*(1) **by** (*unfold-locales*) (*auto simp add: measurable-split-conv*)

lemma (in *nat-filtered-measure*) *predictable-process-const-fun'*[*intro*]:
assumes $f \in \text{borel-measurable } (F \ 0)$
shows *nat-predictable-process* $M \ F \ (\lambda \cdot. f)$
using *assms* **by** (*intro predictable-process-const-fun*[*OF* *measurable-predictable-sigma-snd'*, *THEN* *nat-predictable-process.intro*])

lemma (in *real-filtered-measure*) *predictable-process-const-fun'*[*intro*]:
assumes $f \in \text{borel-measurable } (F \ 0)$
shows *real-predictable-process* $M \ F \ (\lambda \cdot. f)$
using *assms* **by** (*intro predictable-process-const-fun*[*OF* *measurable-predictable-sigma-snd'*, *THEN* *real-predictable-process.intro*])

lemma (in *linearly-filtered-measure*) *predictable-process-const*:
assumes $fst \in \text{borel-measurable } \Sigma_P \ c \in \text{borel-measurable borel}$
shows *predictable-process* $M \ F \ t_0 \ (\lambda i \cdot. c \ i)$
using *assms* **by** (*unfold-locales*) (*simp add: measurable-split-conv*)

lemma (in *linearly-filtered-measure*) *predictable-process-const-const*[*intro*]:
shows *predictable-process* $M \ F \ t_0 \ (\lambda \cdot. \cdot. c)$
by (*unfold-locales*) *simp*

lemma (in *nat-filtered-measure*) *predictable-process-const'*[*intro*]:
assumes $c \in \text{borel-measurable borel}$
shows *nat-predictable-process* $M \ F \ (\lambda i \cdot. c \ i)$
using *assms* **by** (*intro predictable-process-const*[*OF* *measurable-predictable-sigma-fst'*, *THEN* *nat-predictable-process.intro*])

lemma (in *real-filtered-measure*) *predictable-process-const'*[*intro*]:
assumes $c \in \text{borel-measurable borel}$
shows *real-predictable-process* $M \ F \ (\lambda i \cdot. c \ i)$
using *assms* **by** (*intro predictable-process-const*[*OF* *measurable-predictable-sigma-fst'*, *THEN* *real-predictable-process.intro*])

context *predictable-process*
begin

lemma *compose-predictable*:
assumes $fst \in \text{borel-measurable } \Sigma_P \ \text{case-prod } f \in \text{borel-measurable borel}$
shows *predictable-process* $M \ F \ t_0 \ (\lambda i \ \xi. (f \ i) \ (X \ i \ \xi))$
proof
have $(\lambda(i, \xi). (i, X \ i \ \xi)) \in \Sigma_P \rightarrow_M \text{borel} \otimes_M \text{borel}$ **using** *predictable* *assms*(1)

by (auto simp add: measurable-pair-iff measurable-split-conv)
 moreover have $(\lambda(i, \xi). f\ i\ (X\ i\ \xi)) = \text{case-prod } f\ o\ (\lambda(i, \xi). (i, X\ i\ \xi))$ by
fastforce
 ultimately show $(\lambda(i, \xi). f\ i\ (X\ i\ \xi)) \in \text{borel-measurable } \Sigma_P$ **unfolding** *borel-prod*
 using *assms* by *simp*
 qed

lemma *norm-predictable: predictable-process* $M\ F\ t_0\ (\lambda i\ \xi. \text{norm } (X\ i\ \xi))$ **using**
measurable-compose[OF predictable borel-measurable-norm]
 by (unfold-locales) (simp add: prod.case-distrib)

lemma *scaleR-right-predictable:*
 assumes *predictable-process* $M\ F\ t_0\ R$
 shows *predictable-process* $M\ F\ t_0\ (\lambda i\ \xi. (R\ i\ \xi) *_{\mathbb{R}} (X\ i\ \xi))$
 using *predictable predictable-process.predictable[OF assms]* **by** (unfold-locales)
 (auto simp add: measurable-split-conv)

lemma *scaleR-right-const-fun-predictable:*
 assumes $\text{snd} \in \Sigma_P \rightarrow_M F\ t_0\ f \in \text{borel-measurable } (F\ t_0)$
 shows *predictable-process* $M\ F\ t_0\ (\lambda i\ \xi. f\ \xi *_{\mathbb{R}} (X\ i\ \xi))$
 using *assms* **by** (fast intro: *scaleR-right-predictable predictable-process-const-fun*)

lemma *scaleR-right-const-predictable:*
 assumes $\text{fst} \in \text{borel-measurable } \Sigma_P\ c \in \text{borel-measurable borel}$
 shows *predictable-process* $M\ F\ t_0\ (\lambda i\ \xi. c\ i *_{\mathbb{R}} (X\ i\ \xi))$
 using *assms* **by** (fastforce intro: *scaleR-right-predictable predictable-process-const*)

lemma *scaleR-right-const'-predictable: predictable-process* $M\ F\ t_0\ (\lambda i\ \xi. c *_{\mathbb{R}} (X\ i\ \xi))$
 by (fastforce intro: *scaleR-right-predictable*)

lemma *add-predictable:*
 assumes *predictable-process* $M\ F\ t_0\ Y$
 shows *predictable-process* $M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi + Y\ i\ \xi)$
 using *predictable predictable-process.predictable[OF assms]* **by** (unfold-locales)
 (auto simp add: measurable-split-conv)

lemma *diff-predictable:*
 assumes *predictable-process* $M\ F\ t_0\ Y$
 shows *predictable-process* $M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi - Y\ i\ \xi)$
 using *predictable predictable-process.predictable[OF assms]* **by** (unfold-locales)
 (auto simp add: measurable-split-conv)

lemma *uminus-predictable: predictable-process* $M\ F\ t_0\ (-X)$ **using** *scaleR-right-const'-predictable[of -1]* **by** (simp add: fun-Compl-def)

end

Every predictable process is also progressively measurable.

sublocale *predictable-process* \subseteq *progressive-process*

proof (*unfold-locales*)

fix $i :: 'b$ **assume** $asm: t_0 \leq i$

{

fix $S :: ('b \times 'a)$ **set** **assume** $S \in \{\{s <..t\} \times A \mid A \ s \ t. A \in F \ s \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in F \ t_0\}$

hence $(\lambda x. x) - ' S \cap (\{t_0..i\} \times space \ M) \in restrict\text{-}space \ borel \ \{t_0..i\} \otimes_M F$

i

proof

assume $S \in \{\{s <..t\} \times A \mid A \ s \ t. A \in F \ s \wedge t_0 \leq s \wedge s < t\}$

then obtain $s \ t \ A$ **where** $S\text{-is}: S = \{s <..t\} \times A \ t_0 \leq s \ s < t \ A \in F \ s$ **by**

blast

hence $(\lambda x. x) - ' S \cap (\{t_0..i\} \times space \ M) = \{s <..min \ i \ t\} \times A$ **using**

sets.sets-into-space[OF S-is(4)] **by** *auto*

then show $?thesis$ **using** $S\text{-is}$ *sets-F-mono[of s i]* **by** (*cases s ≤ i*) (*fastforce simp add: sets-restrict-space-iff*) +

next

assume $S \in \{\{t_0\} \times A \mid A. A \in F \ t_0\}$

then obtain A **where** $S\text{-is}: S = \{t_0\} \times A \ A \in F \ t_0$ **by** *blast*

hence $(\lambda x. x) - ' S \cap (\{t_0..i\} \times space \ M) = \{t_0\} \times A$ **using** asm *sets.sets-into-space[OF S-is(2)]* **by** *auto*

thus $?thesis$ **using** $S\text{-is}(2)$ *sets-F-mono[OF order-refl asm]* asm **by** (*fastforce simp add: sets-restrict-space-iff*)

qed

hence $(\lambda x. x) - ' S \cap space \ (restrict\text{-}space \ borel \ \{t_0..i\} \otimes_M F \ i) \in restrict\text{-}space \ borel \ \{t_0..i\} \otimes_M F \ i$ **by** (*simp add: space-pair-measure space-F[OF asm]*)

}

moreover have $\{\{s <..t\} \times A \mid A \ s \ t. A \in sets \ (F \ s) \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in sets \ (F \ t_0)\} \subseteq Pow \ (\{t_0..i\} \times space \ M)$ **using** *sets.sets-into-space* **by**

force

ultimately have $(\lambda x. x) \in restrict\text{-}space \ borel \ \{t_0..i\} \otimes_M F \ i \rightarrow_M \Sigma_P$ **using** *space-F[OF asm]* **by** (*intro measurable-sigma-sets[OF sets-predictable-sigma]*)

(*fast, force simp add: space-pair-measure*)

thus case-prod X **is** *borel-measurable* $(restrict\text{-}space \ borel \ \{t_0..i\} \otimes_M F \ i)$ **using**

predictable **by** *simp*

qed

sublocale *nat-predictable-process* \subseteq *nat-progressive-process* ..

sublocale *real-predictable-process* \subseteq *real-progressive-process* ..

The following lemma characterizes predictability in a discrete-time setting.

lemma (*in nat-filtered-measure*) *sets-in-filtration*:

assumes $(\bigcup i. \{i\} \times A \ i) \in \Sigma_P$

shows $A \ (Suc \ i) \in F \ i \ A \ 0 \in F \ 0$

using *assms* **unfolding** *sets-predictable-sigma*

proof (*induction* $(\bigcup i. \{i\} \times A \ i)$ *arbitrary: A*)

case *Basic*

{

assume $\exists S. (\bigcup i. \{i\} \times A \ i) = \{0\} \times S$

then obtain S where $S: (\bigcup i. \{i\} \times A \ i) = \{bot\} \times S$ **unfolding** *bot-nat-def*
 by *blast*
 hence $S \in F \ bot$ **using** *Basic* **by** (*fastforce simp add: times-eq-iff bot-nat-def*)
 moreover have $A \ i = \{\}$ **if** $i \neq bot$ **for** i **using** *that S* **by** *blast*
 moreover have $A \ bot = S$ **using** S **by** *blast*
 ultimately have $A \ (Suc \ i) \in F \ i \ A \ 0 \in F \ 0$ **for** i **unfolding** *bot-nat-def* **by**
 (*auto simp add: bot-nat-def*)
 }
 note $*$ = *this*
 {
 assume $\nexists S. (\bigcup i. \{i\} \times A \ i) = \{0\} \times S$
 then obtain $s \ t \ B$ where $B: (\bigcup i. \{i\} \times A \ i) = \{s<..t\} \times B \ B \in sets \ (F \ s)$
 $s < t$ **using** *Basic* **by** *auto*
 hence $A \ i = B$ **if** $i \in \{s<..t\}$ **for** i **using** *that* **by** *fast*
 moreover have $A \ i = \{\}$ **if** $i \notin \{s<..t\}$ **for** i **using** B **that** **by** *fastforce*
 ultimately have $A \ (Suc \ i) \in F \ i \ A \ 0 \in F \ 0$ **for** i **unfolding** *bot-nat-def* **using**
 $B \ sets-F-mono$ **by** (*auto simp add: bot-nat-def*) (*metis less-Suc-eq-le sets.empty-sets subset-eq*)
 }
 note $**$ = *this*
 show $A \ (Suc \ i) \in sets \ (F \ i) \ A \ 0 \in sets \ (F \ 0)$ **using** $*(1)[of \ i] \ *(2) \ **(1)[of \ i]$
 $**(2)$ **by** *blast+*
 next
 case *Empty*
 {
 case 1
 then show *?case* **using** *Empty* **by** *simp*
 next
 case 2
 then show *?case* **using** *Empty* **by** *simp*
 }
 next
 case (*Compl a*)
 have $a-in: a \subseteq \{0..\} \times space \ M$ **using** *Compl(1) sets.sets-into-space sets-predictable-sigma*
space-predictable-sigma **by** *metis*
 hence $A-in: A \ i \subseteq space \ M$ **for** i **using** *Compl(4)* **by** *blast*
 have $a: a = \{0..\} \times space \ M - (\bigcup i. \{i\} \times A \ i)$ **using** $a-in$ *Compl(4)* **by** *blast*
 also have $\dots = - (\bigcap j. - (\{j\} \times (space \ M - A \ j)))$ **by** *blast*
 also have $\dots = (\bigcup j. \{j\} \times (space \ M - A \ j))$ **by** *blast*
 finally have $*$: $(space \ M - A \ (Suc \ i)) \in F \ i \ (space \ M - A \ 0) \in F \ 0$ **using**
Compl(2,3) **by** *auto*
 {
 case 1
 then show *?case* **using** $*$ $A-in$ **by** (*metis bot-nat-0.extremum double-diff*
sets.Diff sets.top sets-F-mono sets-le-imp-space-le space-F)
 next
 case 2
 then show *?case* **using** $*$ $A-in$ **by** (*metis bot-nat-0.extremum double-diff*
sets.Diff sets.top sets-F-mono sets-le-imp-space-le space-F)

```

}
next
case (Union a)
  have a-in:  $a \ i \subseteq \{0..\} \times \text{space } M$  for  $i$  using Union(1) sets.sets-into-space
sets-predictable-sigma space-predictable-sigma by metis
  hence A-in:  $A \ i \subseteq \text{space } M$  for  $i$  using Union(4) by blast
  have snd  $x \in \text{snd } ' (a \ i \cap (\{fst \ x\} \times \text{space } M))$  if  $x \in a \ i$  for  $i \ x$  using that
a-in by fastforce
  hence a-i:  $a \ i = (\bigcup j. \{j\} \times (\text{snd } ' (a \ i \cap (\{j\} \times \text{space } M))))$  for  $i$  by force
  have A-i:  $A \ i = \text{snd } ' (\bigcup (\text{range } a) \cap (\{i\} \times \text{space } M))$  for  $i$  unfolding
Union(4) using A-in by force
  have *:  $\text{snd } ' (a \ j \cap (\{Suc \ i\} \times \text{space } M)) \in F \ i \ \text{snd } ' (a \ j \cap (\{0\} \times \text{space } M))$ 
 $\in F \ 0$  for  $j$  using Union(2,3)[OF a-i] by auto
  {
    case 1
    have  $(\bigcup j. \text{snd } ' (a \ j \cap (\{Suc \ i\} \times \text{space } M))) \in F \ i$  using * by fast
    moreover have  $(\bigcup j. \text{snd } ' (a \ j \cap (\{Suc \ i\} \times \text{space } M))) = \text{snd } ' (\bigcup (\text{range } a) \cap (\{Suc \ i\} \times \text{space } M))$  by fast
    ultimately show ?case using A-i by metis
  }
  next
  case 2
  have  $(\bigcup j. \text{snd } ' (a \ j \cap (\{0\} \times \text{space } M))) \in F \ 0$  using * by fast
  moreover have  $(\bigcup j. \text{snd } ' (a \ j \cap (\{0\} \times \text{space } M))) = \text{snd } ' (\bigcup (\text{range } a) \cap (\{0\} \times \text{space } M))$  by fast
  ultimately show ?case using A-i by metis
  }
}
qed

```

This leads to the following useful fact.

lemma (in *nat-predictable-process*) *adapted-Suc*: *nat-adapted-process* $M \ F \ (\lambda i. \ X \ (Suc \ i))$

proof (*unfold-locales*, *intro borel-measurableI*)

```

fix S :: 'b set and i assume open-S: open S
have {Suc i} = {i<..Suc i} by fastforce
hence {Suc i}  $\times \text{space } M \in \Sigma_P$  using space-F[symmetric, of i] unfolding
sets-predictable-sigma by (intro sigma-sets.Basic) blast
moreover have case-prod  $X - ' S \cap (UNIV \times \text{space } M) \in \Sigma_P$  unfolding
atLeast-0[symmetric] using open-S by (intro predictableD, simp add: borel-open)
ultimately have case-prod  $X - ' S \cap (\{Suc \ i\} \times \text{space } M) \in \Sigma_P$  unfolding
sets-predictable-sigma using space-F sets.sets-into-space
by (subst Times-Int-distrib1[of {Suc i} UNIV space M, simplified], subst
inf commute, subst Int-assoc[symmetric], subst Int-range-binary)
(intro sigma-sets-Inter binary-in-sigma-sets, fast)+
moreover have case-prod  $X - ' S \cap (\{Suc \ i\} \times \text{space } M) = \{Suc \ i\} \times (X \ (Suc \ i) - ' S \cap \text{space } M)$  by (auto simp add: le-Suc-eq)
moreover have ... =  $(\bigcup j. \{j\} \times (\text{if } j = Suc \ i \text{ then } (X \ (Suc \ i) - ' S \cap \text{space } M) \text{ else } \{\}))$  by (force split: if-splits)
ultimately have  $(\bigcup j. \{j\} \times (\text{if } j = Suc \ i \text{ then } (X \ (Suc \ i) - ' S \cap \text{space } M) \text{ else } \{\})) \in \Sigma_P$  by argo

```

thus $X (Suc\ i) - ' S \cap space\ (F\ i) \in sets\ (F\ i)$ **using** *sets-in-filtration*[*of* $\lambda j.$
if $j = Suc\ i$ *then* $(X\ (Suc\ i) - ' S \cap space\ M)$ *else* $\{\}$] *space-F*[*OF zero-le*] **by**
presburger
qed

The following lemma characterizes predictability in the discrete setting.

theorem *nat-predictable-process-iff*: $nat\text{-}predictable\text{-}process\ M\ F\ X \longleftrightarrow nat\text{-}adapted\text{-}process\ M\ F\ (\lambda i. X\ (Suc\ i)) \wedge X\ 0 \in borel\text{-}measurable\ (F\ 0)$

proof (*intro iffI*)

assume *asm*: $nat\text{-}adapted\text{-}process\ M\ F\ (\lambda i. X\ (Suc\ i)) \wedge X\ 0 \in borel\text{-}measurable\ (F\ 0)$

interpret $nat\text{-}adapted\text{-}process\ M\ F\ \lambda i. X\ (Suc\ i)$ **using** *asm* **by** *blast*

have $(\lambda(x, y). X\ x\ y) \in borel\text{-}measurable\ \Sigma_P$

proof (*intro borel-measurableI*)

fix $S :: 'b\ set$ **assume** *open-S*: $open\ S$

have $\{i\} \times (X\ i - ' S \cap space\ M) \in sets\ \Sigma_P$ **for** i

proof (*cases i*)

case 0

then show *?thesis unfolding sets-predictable-sigma*

using *measurable-sets*[*OF - borel-open*[*OF open-S*], *of* $X\ 0\ F\ 0$] *asm* **by** *auto*

next

case $(Suc\ i)$

have $\{Suc\ i\} = \{i <.. Suc\ i\}$ **by** *fastforce*

then show *?thesis unfolding sets-predictable-sigma*

using *measurable-sets*[*OF adapted borel-open*[*OF open-S*], *of* i]

by (*intro sigma-sets.Basic*, *auto simp add: Suc*)

qed

moreover have $(\lambda(x, y). X\ x\ y) - ' S \cap space\ \Sigma_P = (\bigcup i. \{i\} \times (X\ i - ' S \cap space\ M))$ **by** *fastforce*

ultimately show $(\lambda(x, y). X\ x\ y) - ' S \cap space\ \Sigma_P \in sets\ \Sigma_P$ **by** *simp*

qed

thus $nat\text{-}predictable\text{-}process\ M\ F\ X$ **by** (*unfold-locales*)

next

assume *asm*: $nat\text{-}predictable\text{-}process\ M\ F\ X$

interpret $nat\text{-}predictable\text{-}process\ M\ F\ X$ **by** (*rule asm*)

show $nat\text{-}adapted\text{-}process\ M\ F\ (\lambda i. X\ (Suc\ i)) \wedge X\ 0 \in borel\text{-}measurable\ (F\ 0)$

using *adapted-Suc* **by** *simp*

qed

end

theory *Martingale*

imports *Stochastic-Process Conditional-Expectation-Banach*

begin

7 Martingales

The following locales are necessary for defining martingales.

7.1 Additional Locale Definitions

locale *sigma-finite-adapted-process* = *sigma-finite-filtered-measure* $M\ F\ t_0$ + *adapted-process* $M\ F\ t_0\ X$ **for** $M\ F\ t_0\ X$

locale *nat-sigma-finite-adapted-process* = *sigma-finite-adapted-process* $M\ F\ 0 :: nat\ X$ **for** $M\ F\ X$

locale *real-sigma-finite-adapted-process* = *sigma-finite-adapted-process* $M\ F\ 0 :: real\ X$ **for** $M\ F\ X$

sublocale *nat-sigma-finite-adapted-process* \subseteq *nat-sigma-finite-filtered-measure* ..

sublocale *real-sigma-finite-adapted-process* \subseteq *real-sigma-finite-filtered-measure* ..

locale *finite-adapted-process* = *finite-filtered-measure* $M\ F\ t_0$ + *adapted-process* $M\ F\ t_0\ X$ **for** $M\ F\ t_0\ X$

sublocale *finite-adapted-process* \subseteq *sigma-finite-adapted-process* ..

locale *nat-finite-adapted-process* = *finite-adapted-process* $M\ F\ 0 :: nat\ X$ **for** $M\ F\ X$

locale *real-finite-adapted-process* = *finite-adapted-process* $M\ F\ 0 :: real\ X$ **for** $M\ F\ X$

sublocale *nat-finite-adapted-process* \subseteq *nat-sigma-finite-adapted-process* ..

sublocale *real-finite-adapted-process* \subseteq *real-sigma-finite-adapted-process* ..

locale *sigma-finite-adapted-process-order* = *sigma-finite-adapted-process* $M\ F\ t_0\ X$ **for** $M\ F\ t_0$ **and** $X :: - \Rightarrow - \Rightarrow - :: \{order-topology, ordered-real-vector\}$

locale *nat-sigma-finite-adapted-process-order* = *sigma-finite-adapted-process-order* $M\ F\ 0 :: nat\ X$ **for** $M\ F\ X$

locale *real-sigma-finite-adapted-process-order* = *sigma-finite-adapted-process-order* $M\ F\ 0 :: real\ X$ **for** $M\ F\ X$

sublocale *nat-sigma-finite-adapted-process-order* \subseteq *nat-sigma-finite-adapted-process* ..

sublocale *real-sigma-finite-adapted-process-order* \subseteq *real-sigma-finite-adapted-process* ..

locale *finite-adapted-process-order* = *finite-adapted-process* $M\ F\ t_0\ X$ **for** $M\ F\ t_0$ **and** $X :: - \Rightarrow - \Rightarrow - :: \{order-topology, ordered-real-vector\}$

locale *nat-finite-adapted-process-order* = *finite-adapted-process-order* $M\ F\ 0 :: nat$

```

X for M F X
locale real-finite-adapted-process-order = finite-adapted-process-order M F 0 :: real
X for M F X

sublocale nat-finite-adapted-process-order  $\subseteq$  nat-sigma-finite-adapted-process-order
..
sublocale real-finite-adapted-process-order  $\subseteq$  real-sigma-finite-adapted-process-order
..

locale sigma-finite-adapted-process-linorder = sigma-finite-adapted-process-order
M F t0 X for M F t0 and X :: -  $\Rightarrow$  -  $\Rightarrow$  - :: {linorder-topology}

locale nat-sigma-finite-adapted-process-linorder = sigma-finite-adapted-process-linorder
M F 0 :: nat X for M F X
locale real-sigma-finite-adapted-process-linorder = sigma-finite-adapted-process-linorder
M F 0 :: real X for M F X

sublocale nat-sigma-finite-adapted-process-linorder  $\subseteq$  nat-sigma-finite-adapted-process-order
..
sublocale real-sigma-finite-adapted-process-linorder  $\subseteq$  real-sigma-finite-adapted-process-order
..

locale finite-adapted-process-linorder = finite-adapted-process-order M F t0 X for
M F t0 and X :: -  $\Rightarrow$  -  $\Rightarrow$  - :: {linorder-topology}

locale nat-finite-adapted-process-linorder = finite-adapted-process-linorder M F 0
:: nat X for M F X
locale real-finite-adapted-process-linorder = finite-adapted-process-linorder M F 0
:: real X for M F X

sublocale nat-finite-adapted-process-linorder  $\subseteq$  nat-sigma-finite-adapted-process-linorder
..
sublocale real-finite-adapted-process-linorder  $\subseteq$  real-sigma-finite-adapted-process-linorder
..

```

7.2 Martingale

A martingale is an adapted process where the expected value of the next observation, given all past observations, is equal to the current value.

```

locale martingale = sigma-finite-adapted-process +
  assumes integrable:  $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$ 
  and martingale-property:  $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE \xi \text{ in } M. X i \xi =$ 
cond-exp M (F i) (X j) \xi

```

```

locale martingale-order = martingale M F t0 X for M F t0 and X :: -  $\Rightarrow$  -  $\Rightarrow$  -
:: {order-topology, ordered-real-vector}
locale martingale-linorder = martingale M F t0 X for M F t0 and X :: -  $\Rightarrow$  -  $\Rightarrow$ 

```


- :: {linorder-topology, ordered-real-vector}
sublocale martingale-linorder \subseteq martingale-order ..

lemma (in sigma-finite-filtered-measure) martingale-const-fun[intro]:
 assumes integrable M f $f \in$ borel-measurable $(F\ t_0)$
 shows martingale M $F\ t_0$ $(\lambda\cdot. f)$
 using assms sigma-finite-subalgebra.cond-exp-F-meas[OF - assms(1), THEN AE-symmetric]
 borel-measurable-mono
 by (unfold-locale) blast+

lemma (in sigma-finite-filtered-measure) martingale-cond-exp[intro]:
 assumes integrable M f
 shows martingale M $F\ t_0$ $(\lambda i. \text{cond-exp } M\ (F\ i)\ f)$
 using sigma-finite-subalgebra.borel-measurable-cond-exp' borel-measurable-cond-exp

 by (unfold-locale) (auto intro: sigma-finite-subalgebra.cond-exp-nested-subalg[OF - assms] simp add: subalgebra-F subalgebras)

corollary (in sigma-finite-filtered-measure) martingale-zero[intro]: martingale M $F\ t_0$ $(\lambda\cdot. 0)$ by fastforce

corollary (in finite-filtered-measure) martingale-const[intro]: martingale M $F\ t_0$ $(\lambda\cdot. c)$ by fastforce

7.3 Submartingale

A submartingale is an adapted process where, the expected value of the next observation, given all past observations, is greater than or equal to the current value.

locale submartingale = sigma-finite-adapted-process-order +
 assumes integrable: $\bigwedge i. t_0 \leq i \implies \text{integrable } M\ (X\ i)$
 and submartingale-property: $\bigwedge i\ j. t_0 \leq i \implies i \leq j \implies \text{AE } \xi \text{ in } M. X\ i\ \xi \leq \text{cond-exp } M\ (F\ i)\ (X\ j)\ \xi$

locale submartingale-linorder = submartingale $M\ F\ t_0\ X$ for $M\ F\ t_0$ and X :: -
 \Rightarrow - \Rightarrow - :: {linorder-topology}

sublocale martingale-order \subseteq submartingale using martingale-property by (unfold-locale)
 (force simp add: integrable)+
sublocale martingale-linorder \subseteq submartingale-linorder ..

7.4 Supermartingale

A supermartingale is an adapted process where, the expected value of the next observation, given all past observations, is less than or equal to the current value.

locale supermartingale = sigma-finite-adapted-process-order +
 assumes integrable: $\bigwedge i. t_0 \leq i \implies \text{integrable } M\ (X\ i)$

and *supermartingale-property*: $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE \xi \text{ in } M. X i \xi \geq \text{cond-exp } M (F i) (X j) \xi$

locale *supermartingale-linorder* = *supermartingale* $M F t_0 X$ **for** $M F t_0$ **and** X
 $:: - \Rightarrow - \Rightarrow - :: \{\text{linorder-topology}\}$

sublocale *martingale-order* \subseteq *supermartingale* **using** *martingale-property* **by** (*unfold-locales*)
(*force simp add: integrable*)
sublocale *martingale-linorder* \subseteq *supermartingale-linorder* ..

A stochastic process is a martingale, if and only if it is both a submartingale and a supermartingale.

lemma *martingale-iff*:

shows *martingale* $M F t_0 X \longleftrightarrow \text{submartingale } M F t_0 X \wedge \text{supermartingale } M F t_0 X$

proof (*rule iffI*)

assume *asm*: *martingale* $M F t_0 X$

interpret *martingale-order* $M F t_0 X$ **by** (*intro martingale-order.intro asm*)

show *submartingale* $M F t_0 X \wedge \text{supermartingale } M F t_0 X$ **using** *submartingale-axioms supermartingale-axioms* **by** *blast*

next

assume *asm*: *submartingale* $M F t_0 X \wedge \text{supermartingale } M F t_0 X$

interpret *submartingale* $M F t_0 X$ **by** (*simp add: asm*)

interpret *supermartingale* $M F t_0 X$ **by** (*simp add: asm*)

show *martingale* $M F t_0 X$ **using** *submartingale-property supermartingale-property* **by** (*unfold-locales*) (*intro integrable, blast, force*)

qed

7.5 Martingale Lemmas

In the following segment, we cover basic properties of martingales.

context *martingale*

begin

lemma *cond-exp-diff-eq-zero*:

assumes $t_0 \leq i \leq j$

shows $AE \xi \text{ in } M. \text{cond-exp } M (F i) (\lambda \xi. X j \xi - X i \xi) \xi = 0$

using *martingale-property*[*OF assms*] *assms*

sigma-finite-subalgebra.cond-exp-F-meas[*OF - integrable adapted, of i*]

sigma-finite-subalgebra.cond-exp-diff[*OF - integrable(1,1), of F i j i*] **by**

fastforce

lemma *set-integral-eq*:

assumes $A \in F i t_0 \leq i \leq j$

shows *set-lebesgue-integral* $M A (X i) = \text{set-lebesgue-integral } M A (X j)$

proof –

interpret *sigma-finite-subalgebra* $M F i$ **using** *assms(2)* **by** *blast*

have $\int x \in A. X i x \partial M = \int x \in A. \text{cond-exp } M (F i) (X j) x \partial M$ **using** *martingale-property*[*OF assms(2,3)*] *borel-measurable-cond-exp'* *assms subalgebras*

subalgebra-def **by** (intro set-lebesgue-integral-cong-AE[OF - random-variable]) fast-force+
also have ... = $\int x \in A. X j x \partial M$ **using** *assms* **by** (auto simp: integrable intro: cond-exp-set-integral[symmetric])
finally show ?thesis .
qed

lemma scaleR-const[intro]:
shows martingale $M F t_0 (\lambda i x. c *_R X i x)$
proof –
{
 fix $i j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
 interpret sigma-finite-subalgebra $M F i$ **using** *asm* **by** blast
 have AE x in $M. c *_R X i x = \text{cond-exp } M (F i) (\lambda x. c *_R X j x) x$ **using** *asm* cond-exp-scaleR-right[OF integrable, of j , THEN AE-symmetric] martingale-property[OF *asm*] **by** force
}
thus ?thesis **by** (unfold-locales) (auto simp add: integrable martingale.integrable)
qed

lemma uminus[intro]:
shows martingale $M F t_0 (- X)$
using scaleR-const[of -1] **by** (force intro: back-subst[of martingale $M F t_0$])

lemma add[intro]:
assumes martingale $M F t_0 Y$
shows martingale $M F t_0 (\lambda i \xi. X i \xi + Y i \xi)$
proof –
interpret Y : martingale $M F t_0 Y$ **by** (rule *assms*)
{
 fix $i j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
 hence AE ξ in $M. X i \xi + Y i \xi = \text{cond-exp } M (F i) (\lambda x. X j x + Y j x) \xi$
 using sigma-finite-subalgebra.cond-exp-add[OF - integrable martingale.integrable[OF *assms*], of $F i j j$, THEN AE-symmetric]
 martingale-property[OF *asm*] martingale.martingale-property[OF *assms* *asm*] **by** force
}
thus ?thesis **using** *assms*
by (unfold-locales) (auto simp add: integrable martingale.integrable)
qed

lemma diff[intro]:
assumes martingale $M F t_0 Y$
shows martingale $M F t_0 (\lambda i x. X i x - Y i x)$
proof –
interpret Y : martingale $M F t_0 Y$ **by** (rule *assms*)
{
 fix $i j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
 hence AE ξ in $M. X i \xi - Y i \xi = \text{cond-exp } M (F i) (\lambda x. X j x - Y j x) \xi$

```

    using sigma-finite-subalgebra.cond-exp-diff[OF - integrable martingale.integrable[OF
assms], of F i j j, THEN AE-symmetric]
      martingale-property[OF asm] martingale.martingale-property[OF assms
asm] by fastforce
  }
  thus ?thesis using assms by (unfold-locales) (auto simp add: integrable martin-
gale.integrable)
qed

end

```

Using properties of the conditional expectation, we present the following alternative characterizations of martingales.

```

lemma (in sigma-finite-adapted-process) martingale-of-cond-exp-diff-eq-zero:
  assumes integrable:  $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$ 
    and diff-zero:  $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE \ x \text{ in } M. \text{ cond-exp } M (F i) (\lambda \xi. X j \xi - X i \xi) \ x = 0$ 
  shows martingale M F t0 X
proof
  {
    fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
    thus AE  $\xi$  in M.  $X i \xi = \text{cond-exp } M (F i) (X j) \xi$ 
    using diff-zero[OF asm] sigma-finite-subalgebra.cond-exp-diff[OF - integrable(1,1),
of F i j i]
      sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] by
fastforce
  }
qed (intro integrable)

```

```

lemma (in sigma-finite-adapted-process) martingale-of-set-integral-eq:
  assumes integrable:  $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$ 
    and  $\bigwedge A \ i \ j. t_0 \leq i \implies i \leq j \implies A \in F i \implies \text{set-lebesgue-integral } M A (X i) = \text{set-lebesgue-integral } M A (X j)$ 
  shows martingale M F t0 X
proof (unfold-locales)
  fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
  interpret sigma-finite-subalgebra M F i using asm by blast
  interpret r: sigma-finite-measure restr-to-subalg M (F i) by (simp add: sigma-fin-subalg)
  {
    fix A assume A  $\in \text{restr-to-subalg } M (F i)$ 
    hence *:  $A \in F i$  using sets-restr-to-subalg subalgebras asm by blast
    have set-lebesgue-integral (restr-to-subalg M (F i)) A (X i) = set-lebesgue-integral
M A (X i) using * subalg asm by (auto simp: set-lebesgue-integral-def intro: inte-
gral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator)
    also have ... = set-lebesgue-integral M A (cond-exp M (F i) (X j)) using *
assms(2)[OF asm] cond-exp-set-integral[OF integrable] asm by auto
    finally have set-lebesgue-integral (restr-to-subalg M (F i)) A (X i) = set-lebesgue-integral
(restr-to-subalg M (F i)) A (cond-exp M (F i) (X j)) using * subalg by (auto simp:
set-lebesgue-integral-def intro!: integral-subalgebra2[symmetric] borel-measurable-scaleR

```

borel-measurable-cond-exp borel-measurable-indicator)
 }
 hence $AE \xi$ in $restr\text{-}to\text{-}subalg\ M\ (F\ i)$. $X\ i\ \xi = cond\text{-}exp\ M\ (F\ i)\ (X\ j)\ \xi$ **using** asm **by** (intro $r.density\text{-}unique\text{-}banach$, auto intro: $integrable\text{-}in\text{-}subalg\ subalg\ borel\text{-}measurable\text{-}cond\text{-}exp\ integrable$)
 thus $AE \xi$ in M . $X\ i\ \xi = cond\text{-}exp\ M\ (F\ i)\ (X\ j)\ \xi$ **using** $AE\text{-}restr\text{-}to\text{-}subalg\ [OF\ subalg]$ **by** blast
 qed (simp add: $integrable$)

7.6 Submartingale Lemmas

context $submartingale$
 begin

lemma $cond\text{-}exp\text{-}diff\text{-}nonneg$:
 assumes $t_0 \leq i\ i \leq j$
 shows $AE\ x$ in M . $cond\text{-}exp\ M\ (F\ i)\ (\lambda\xi. X\ j\ \xi - X\ i\ \xi)\ x \geq 0$
using $submartingale\text{-}property\ [OF\ assms]$ $assms\ sigma\text{-}finite\text{-}subalgebra.cond\text{-}exp\text{-}diff\ [OF\ -\ integrable(1,1),\ of\ -\ j\ i]\ sigma\text{-}finite\text{-}subalgebra.cond\text{-}exp\text{-}F\text{-}meas\ [OF\ -\ integrable\ adapted,\ of\ i]$ **by** fastforce

lemma $add\ [intro]$:
 assumes $submartingale\ M\ F\ t_0\ Y$
 shows $submartingale\ M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi + Y\ i\ \xi)$
proof –
 interpret Y : $submartingale\ M\ F\ t_0\ Y$ **by** (rule $assms$)
 {
 fix $i\ j :: 'b$ **assume** $asm: t_0 \leq i\ i \leq j$
 hence $AE\ \xi$ in M . $X\ i\ \xi + Y\ i\ \xi \leq cond\text{-}exp\ M\ (F\ i)\ (\lambda x. X\ j\ x + Y\ j\ x)\ \xi$
 using $sigma\text{-}finite\text{-}subalgebra.cond\text{-}exp\text{-}add\ [OF\ -\ integrable\ submartingale.integrable\ [OF\ assms],\ of\ F\ i\ j\ j]$
 $submartingale\text{-}property\ [OF\ asm]\ submartingale.submartingale\text{-}property\ [OF\ assms\ asm]\ add\text{-}mono\ [of\ X\ i\ -\ Y\ i\ -]$ **by** force
 }
 thus ?thesis **using** $assms$ **by** (unfold-locale) (auto simp add: $borel\text{-}measurable\text{-}add\ random\text{-}variable\ adapted\ integrable\ Y.random\text{-}variable\ Y.adapted\ submartingale.integrable$)

qed

lemma $diff\ [intro]$:
 assumes $supermartingale\ M\ F\ t_0\ Y$
 shows $supermartingale\ M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi - Y\ i\ \xi)$
proof –
 interpret Y : $supermartingale\ M\ F\ t_0\ Y$ **by** (rule $assms$)
 {
 fix $i\ j :: 'b$ **assume** $asm: t_0 \leq i\ i \leq j$
 hence $AE\ \xi$ in M . $X\ i\ \xi - Y\ i\ \xi \leq cond\text{-}exp\ M\ (F\ i)\ (\lambda x. X\ j\ x - Y\ j\ x)\ \xi$
 using $sigma\text{-}finite\text{-}subalgebra.cond\text{-}exp\text{-}diff\ [OF\ -\ integrable\ supermartingale.integrable\ [OF\ assms],\ of\ F\ i\ j\ j]$

```

      submartingale-property[OF asm] supermartingale.supermartingale-property[OF
assms asm] diff-mono[of  $X$   $i$  - - -  $Y$   $i$  -] by force
    }
    thus ?thesis using assms by (unfold-locale) (auto simp add: borel-measurable-diff
random-variable adapted integrable  $Y$ .random-variable  $Y$ .adapted supermartingale.integrable)

qed

```

```

lemma scaleR-nonneg:
  assumes  $c \geq 0$ 
  shows submartingale  $M$   $F$   $t_0$  ( $\lambda i \xi. c *_{\mathbb{R}} X i \xi$ )
proof
  {
    fix  $i j :: 'b$  assume asm:  $t_0 \leq i \leq j$ 
    thus  $\mathbb{E} \xi$  in  $M. c *_{\mathbb{R}} X i \xi \leq \text{cond-exp } M (F i) (\lambda \xi. c *_{\mathbb{R}} X j \xi) \xi$ 
      using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of  $F i$ 
 $j$   $c$ ] submartingale-property[OF asm] by (fastforce intro!: scaleR-left-mono[OF -
assms])
  }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)

```

```

lemma scaleR-le-zero:
  assumes  $c \leq 0$ 
  shows supermartingale  $M$   $F$   $t_0$  ( $\lambda i \xi. c *_{\mathbb{R}} X i \xi$ )
proof
  {
    fix  $i j :: 'b$  assume asm:  $t_0 \leq i \leq j$ 
    thus  $\mathbb{E} \xi$  in  $M. c *_{\mathbb{R}} X i \xi \geq \text{cond-exp } M (F i) (\lambda \xi. c *_{\mathbb{R}} X j \xi) \xi$ 
      using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of  $F i$   $j$ 
 $c$ ] submartingale-property[OF asm]
      by (fastforce intro!: scaleR-left-mono-neg[OF - assms])
  }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)

```

```

lemma uminus[intro]:
  shows supermartingale  $M$   $F$   $t_0$  ( $- X$ )
  unfolding fun-Comp-def using scaleR-le-zero[of  $-1$ ] by simp

end

```

```

context submartingale-linorder
begin

```

```

lemma set-integral-le:
  assumes  $A \in F$   $i t_0 \leq i \leq j$ 
  shows set-lebesgue-integral  $M$   $A$  ( $X i$ )  $\leq$  set-lebesgue-integral  $M$   $A$  ( $X j$ )
  using submartingale-property[OF assms(2), of  $j$ ] assms subalgebras

```

by (*subst sigma-finite-subalgebra.cond-exp-set-integral*[*OF* - *integrable assms*(1), of *j*])
 (*auto intro!*: *scaleR-left-mono integral-mono-AE-banach integrable-mult-indicator integrable simp add: subalgebra-def set-lebesgue-integral-def*)

lemma *max*:

assumes *submartingale-linorder* *M F t₀ Y*
shows *submartingale-linorder* *M F t₀ (λ i ξ. max (X i ξ) (Y i ξ))*
proof (*unfold-locales*)
interpret *Y*: *submartingale-linorder* *M F t₀ Y* **by** (*rule assms*)
 {
fix *i j* :: 'b **assume** *asm*: *t₀ ≤ i i ≤ j*
have *AE ξ in M. max (X i ξ) (Y i ξ) ≤ max (cond-exp M (F i) (X j) ξ)*
(cond-exp M (F i) (Y j) ξ) **using** *submartingale-property Y.submartingale-property*
asm unfolding max-def by fastforce
thus *AE ξ in M. max (X i ξ) (Y i ξ) ≤ cond-exp M (F i) (λ ξ. max (X j ξ) (Y*
j ξ)) ξ **using** *sigma-finite-subalgebra.cond-exp-max*[*OF* - *integrable Y.integrable, of*
F i j j] *asm by (fast intro: order.trans)*
 }
show $\bigwedge i. t_0 \leq i \implies (\lambda \xi. \max (X i \xi) (Y i \xi)) \in \text{borel-measurable } (F i) \bigwedge i.$
 $t_0 \leq i \implies \text{integrable } M (\lambda \xi. \max (X i \xi) (Y i \xi))$ **by** (*force intro: Y.integrable*
integrable assms)
qed

lemma *max-0*:

shows *submartingale-linorder* *M F t₀ (λ i ξ. max 0 (X i ξ))*
proof –
interpret *zero*: *martingale-linorder* *M F t₀ λ- -. 0* **by** (*force intro: martin-*
gale-linorder.intro martingale-order.intro)
show *?thesis* **by** (*intro zero.max submartingale-linorder.intro submartingale-axioms*)
qed

end

lemma (*in sigma-finite-adapted-process-order*) *submartingale-of-cond-exp-diff-nonneg*:

assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
and *diff-nonneg*: $\bigwedge i j. t_0 \leq i \implies i \leq j \implies \text{AE } x \text{ in } M. \text{cond-exp } M (F i)$
 $(\lambda \xi. X j \xi - X i \xi) x \geq 0$
shows *submartingale* *M F t₀ X*
proof (*unfold-locales*)
 {
fix *i j* :: 'b **assume** *asm*: *t₀ ≤ i i ≤ j*
thus *AE ξ in M. X i ξ ≤ cond-exp M (F i) (X j) ξ*
using *diff-nonneg*[*OF asm*] *sigma-finite-subalgebra.cond-exp-diff*[*OF* - *inte-*
grable(1,1), of F i j i]
sigma-finite-subalgebra.cond-exp-F-meas[*OF* - *integrable adapted, of i*] **by**
fastforce
 }
qed (*intro integrable*)

```

lemma (in sigma-finite-adapted-process-linorder) submartingale-of-set-integral-le:
  assumes integrable:  $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$ 
    and  $\bigwedge A i j. t_0 \leq i \implies i \leq j \implies A \in F i \implies \text{set-lebesgue-integral } M A (X i) \leq \text{set-lebesgue-integral } M A (X j)$ 
  shows submartingale  $M F t_0 X$ 
proof (unfold-locales)
  {
    fix  $i j :: 'b$  assume  $asm: t_0 \leq i \leq j$ 
    interpret  $r: \text{sigma-finite-measure restr-to-subalg } M (F i)$  using  $asm \text{ sigma-finite-subalgebra.sigma-fin-subalg}$ 
by blast
    {
      fix  $A$  assume  $A \in \text{restr-to-subalg } M (F i)$ 
      hence  $*$ :  $A \in F i$  using  $asm \text{ sets-restr-to-subalg subalgebras}$  by blast
      have  $\text{set-lebesgue-integral } (\text{restr-to-subalg } M (F i)) A (X i) = \text{set-lebesgue-integral } M A (X i)$  using  $*$   $asm \text{ subalgebras}$  by (auto simp: set-lebesgue-integral-def intro: integral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator)
      also have  $\dots \leq \text{set-lebesgue-integral } M A (\text{cond-exp } M (F i) (X j))$  using
       $*$  assms(2)[OF asm]  $asm \text{ sigma-finite-subalgebra.cond-exp-set-integral[OF - integrable]}$  by fastforce
      also have  $\dots = \text{set-lebesgue-integral } (\text{restr-to-subalg } M (F i)) A (\text{cond-exp } M (F i) (X j))$  using  $*$   $asm \text{ subalgebras}$  by (auto simp: set-lebesgue-integral-def intro!: integral-subalgebra2[symmetric] borel-measurable-scaleR borel-measurable-cond-exp borel-measurable-indicator)
      finally have  $0 \leq \text{set-lebesgue-integral } (\text{restr-to-subalg } M (F i)) A (\lambda \xi. \text{cond-exp } M (F i) (X j) \xi - X i \xi)$  using  $*$   $asm \text{ subalgebras}$  by (subst set-integral-diff, auto simp add: set-integrable-def sets-restr-to-subalg intro!: integrable adapted integrable-in-subalg borel-measurable-scaleR borel-measurable-indicator borel-measurable-cond-exp integrable-mult-indicator)
    }
    hence  $AE \xi \text{ in } \text{restr-to-subalg } M (F i). 0 \leq \text{cond-exp } M (F i) (X j) \xi - X i \xi$ 
    by (intro  $r.\text{density-nonneg integrable-in-subalg}$   $asm \text{ subalgebras borel-measurable-diff borel-measurable-cond-exp adapted Bochner-Integration.integrable-diff integrable-cond-exp integrable}$ )
    thus  $AE \xi \text{ in } M. X i \xi \leq \text{cond-exp } M (F i) (X j) \xi$  using  $AE\text{-restr-to-subalg[OF subalgebras]}$   $asm$  by simp
  }
qed (intro integrable)

```

7.7 Supermartingale Lemmas

The following lemmas are exact duals of the ones for submartingales.

```

context supermartingale
begin

```

```

lemma cond-exp-diff-nonneg:
  assumes  $t_0 \leq i \leq j$ 
  shows  $AE x \text{ in } M. \text{cond-exp } M (F i) (\lambda \xi. X i \xi - X j \xi) x \geq 0$ 
  using assms supermartingale-property[OF assms] sigma-finite-subalgebra.cond-exp-diff[OF

```


- *integrable*(1,1), of *F i i j*]
 sigma-finite-subalgebra.cond-exp-F-meas[*OF - integrable adapted*, of *i*] **by**
fastforce

lemma *add*[*intro*]:
 assumes *supermartingale M F t₀ Y*
 shows *supermartingale M F t₀ (λ i ξ. X i ξ + Y i ξ)*
proof -
 interpret *Y: supermartingale M F t₀ Y* **by** (*rule assms*)
 {
 fix *i j :: 'b* **assume** *asm: t₀ ≤ i i ≤ j*
 hence *AE ξ in M. X i ξ + Y i ξ ≥ cond-exp M (F i) (λ x. X j x + Y j x) ξ*
 using *sigma-finite-subalgebra.cond-exp-add*[*OF - integrable supermartingale.integrable*[*OF*
assms], of *F i j j*]
 supermartingale-property[*OF asm*] *supermartingale.supermartingale-property*[*OF*
assms asm] *add-mono*[of - *X i - - Y i -*] **by** *force*
 }
 thus *?thesis using assms by (unfold-locale) (auto simp add: borel-measurable-add*
random-variable adapted integrable Y.random-variable Y.adapted supermartingale.integrable)

qed

lemma *diff*[*intro*]:
 assumes *submartingale M F t₀ Y*
 shows *supermartingale M F t₀ (λ i ξ. X i ξ - Y i ξ)*
proof -
 interpret *Y: submartingale M F t₀ Y* **by** (*rule assms*)
 {
 fix *i j :: 'b* **assume** *asm: t₀ ≤ i i ≤ j*
 hence *AE ξ in M. X i ξ - Y i ξ ≥ cond-exp M (F i) (λ x. X j x - Y j x) ξ*
 using *sigma-finite-subalgebra.cond-exp-diff*[*OF - integrable submartingale.integrable*[*OF*
assms], of *F i j j*, *unfolded fun-diff-def*]
 supermartingale-property[*OF asm*] *submartingale.submartingale-property*[*OF*
assms asm] *diff-mono*[of - *X i - Y i -*] **by** *force*
 }
 thus *?thesis using assms by (unfold-locale) (auto simp add: borel-measurable-diff*
random-variable adapted integrable Y.random-variable Y.adapted submartingale.integrable)

qed

lemma *scaleR-nonneg*:
 assumes *c ≥ 0*
 shows *supermartingale M F t₀ (λ i ξ. c *_R X i ξ)*
proof
 {
 fix *i j :: 'b* **assume** *asm: t₀ ≤ i i ≤ j*
 thus *AE ξ in M. c *_R X i ξ ≥ cond-exp M (F i) (λ ξ. c *_R X j ξ) ξ*
 using *sigma-finite-subalgebra.cond-exp-scaleR-right*[*OF - integrable*, of *F i*
j c] *supermartingale-property*[*OF asm*] **by** (*fastforce intro! scaleR-left-mono*[*OF -*

```

assms])
}
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)

lemma scaleR-le-zero:
  assumes  $c \leq 0$ 
  shows submartingale  $M F t_0 (\lambda i \xi. c *_{\mathbb{R}} X i \xi)$ 
proof
  {
    fix  $i j :: 'b$  assume asm:  $t_0 \leq i \leq j$ 
    thus  $AE \xi \text{ in } M. c *_{\mathbb{R}} X i \xi \leq \text{cond-exp } M (F i) (\lambda \xi. c *_{\mathbb{R}} X j \xi) \xi$ 
      using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of  $F i j c$ ]
      supermartingale-property[OF asm] by (fastforce intro!: scaleR-left-mono-neg[OF -
      assms])
  }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)

lemma uminus[intro]:
  shows submartingale  $M F t_0 (- X)$ 
  unfolding fun-Comp-def using scaleR-le-zero[of  $-1$ ] by simp

end

context supermartingale-linorder
begin

lemma set-integral-ge:
  assumes  $A \in F i t_0 \leq i \leq j$ 
  shows set-lebesgue-integral  $M A (X i) \geq \text{set-lebesgue-integral } M A (X j)$ 
  using supermartingale-property[OF assms(2), of  $j$ ] assms subalgebras
  by (subst sigma-finite-subalgebra.cond-exp-set-integral[OF - integrable assms(1),
  of  $j$ ])
  (auto intro!: scaleR-left-mono integral-mono-AE-banach integrable-mult-indicator
  integrable simp add: subalgebra-def set-lebesgue-integral-def)

lemma min:
  assumes supermartingale-linorder  $M F t_0 Y$ 
  shows supermartingale-linorder  $M F t_0 (\lambda i \xi. \min (X i \xi) (Y i \xi))$ 
proof (unfold-locales)
  interpret  $Y: \text{supermartingale-linorder } M F t_0 Y$  by (rule assms)
  {
    fix  $i j :: 'b$  assume asm:  $t_0 \leq i \leq j$ 
    have  $AE \xi \text{ in } M. \min (X i \xi) (Y i \xi) \geq \min (\text{cond-exp } M (F i) (X j) \xi) (\text{cond-exp } M (F i) (Y j) \xi)$ 
      using supermartingale-property  $Y.$ supermartingale-property asm
      unfolding min-def by fastforce
    thus  $AE \xi \text{ in } M. \min (X i \xi) (Y i \xi) \geq \text{cond-exp } M (F i) (\lambda \xi. \min (X j \xi) (Y j \xi)) \xi$ 
      using sigma-finite-subalgebra.cond-exp-min[OF - integrable  $Y.$ integrable, of

```

$F\ i\ j\ j]$ *asm* **by** (*fast intro: order.trans*)
}
show $\bigwedge i. t_0 \leq i \implies (\lambda \xi. \min (X\ i\ \xi) (Y\ i\ \xi)) \in \text{borel-measurable } (F\ i) \bigwedge i.$
 $t_0 \leq i \implies \text{integrable } M\ (\lambda \xi. \min (X\ i\ \xi) (Y\ i\ \xi))$ **by** (*force intro: Y.integrable*
integrable assms)
qed

lemma *min-0*:

shows *supermartingale-linorder* $M\ F\ t_0\ (\lambda i\ \xi. \min\ 0\ (X\ i\ \xi))$
proof –
interpret *zero: martingale-linorder* $M\ F\ t_0\ \lambda -.\ 0$ **by** (*force intro: martin-*
gale-linorder.intro)
show *?thesis* **by** (*intro zero.min supermartingale-linorder.intro supermartin-*
gale-axioms)
qed

end

lemma (*in sigma-finite-adapted-process-order*) *supermartingale-of-cond-exp-diff-le-zero*:

assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M\ (X\ i)$
and *diff-le-zero*: $\bigwedge i\ j. t_0 \leq i \implies i \leq j \implies \text{AE } x \text{ in } M. \text{ cond-exp } M\ (F\ i)$
 $(\lambda \xi. X\ j\ \xi - X\ i\ \xi)\ x \leq 0$
shows *supermartingale* $M\ F\ t_0\ X$
proof
{
fix $i\ j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
thus *AE* ξ *in* $M. X\ i\ \xi \geq \text{cond-exp } M\ (F\ i)\ (X\ j)\ \xi$
using *diff-le-zero*[*OF asm*] *sigma-finite-subalgebra.cond-exp-diff*[*OF - inte-*
grable(1,1), of F i j i]
sigma-finite-subalgebra.cond-exp-F-meas[*OF - integrable adapted, of i*] **by**
fastforce
}
qed (*intro integrable*)

lemma (*in sigma-finite-adapted-process-linorder*) *supermartingale-of-set-integral-ge*:

assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M\ (X\ i)$
and $\bigwedge A\ i\ j. t_0 \leq i \implies i \leq j \implies A \in F\ i \implies \text{set-lebesgue-integral } M\ A\ (X\ j) \leq \text{set-lebesgue-integral } M\ A\ (X\ i)$
shows *supermartingale* $M\ F\ t_0\ X$
proof –
interpret *-: adapted-process* $M\ F\ t_0\ -X$ **by** (*rule uminus-adapted*)
interpret *uminus-X*: *sigma-finite-adapted-process-linorder* $M\ F\ t_0\ -X$..
note $*$ = *set-integral-uminus*[*unfolded set-integrable-def, OF integrable-mult-indicator*[*OF*
- integrable]]
have *supermartingale* $M\ F\ t_0\ (-(-X))$
using *ord-eq-le-trans*[*OF * ord-le-eq-trans*[*OF le-imp-neg-le*[*OF assms(2)*]] **[symmetric]*]]
subalgebras
by (*intro submartingale.uminus uminus-X.submartingale-of-set-integral-le*)
(clarsimp simp add: fun-Compl-def subalgebra-def integrable | fastforce)+

thus *?thesis* **unfolding** *fun-Compl-def* **by** *simp*
qed

Many of the statements we have made concerning martingales can be simplified when the indexing set is the natural numbers. Given a point in time $i \in \mathbb{N}$, it suffices to consider the successor $i + (1::'a)$, instead of all future times $i \leq j$.

7.8 Discrete Time Martingales

locale *nat-martingale* = *martingale* $M\ F\ 0 :: nat\ X$ **for** $M\ F\ X$

locale *nat-submartingale* = *submartingale* $M\ F\ 0 :: nat\ X$ **for** $M\ F\ X$

locale *nat-supermartingale* = *supermartingale* $M\ F\ 0 :: nat\ X$ **for** $M\ F\ X$

locale *nat-submartingale-linorder* = *submartingale-linorder* $M\ F\ 0 :: nat\ X$ **for** $M\ F\ X$

locale *nat-supermartingale-linorder* = *supermartingale-linorder* $M\ F\ 0 :: nat\ X$ **for** $M\ F\ X$

sublocale *nat-submartingale-linorder* \subseteq *nat-submartingale* ..

sublocale *nat-supermartingale-linorder* \subseteq *nat-supermartingale* ..

A predictable martingale is necessarily constant.

lemma (**in** *nat-martingale*) *predictable-const*:

assumes *nat-predictable-process* $M\ F\ X$

shows $AE\ \xi\ in\ M. X\ i\ \xi = X\ j\ \xi$

proof –

have *: $AE\ \xi\ in\ M. X\ i\ \xi = X\ 0\ \xi$ **for** i

proof (*induction* i)

case 0

then show *?case* **by** (*simp* *add*: *bot-nat-def*)

next

case (*Suc* i)

interpret S : *nat-adapted-process* $M\ F\ \lambda i. X\ (Suc\ i)$ **by** (*intro* *nat-predictable-process.adapted-Suc* *assms*)

show *?case* **using** *Suc* $S.adapted[of\ i]$ *martingale-property*[*OF* - *le-SucI*, *of* i]
sigma-finite-subalgebra.cond-exp-F-meas[*OF* - *integrable*, *of* $F\ i\ Suc\ i$] **by** *fastforce*

qed

show *?thesis* **using** *[*of* i] *[*of* j] **by** *force*

qed

lemma (**in** *nat-sigma-finite-adapted-process*) *martingale-of-set-integral-eq-Suc*:

assumes *integrable*: $\bigwedge i. integrable\ M\ (X\ i)$

and $\bigwedge A\ i. A \in F\ i \implies set-lebesgue-integral\ M\ A\ (X\ i) = set-lebesgue-integral\ M\ A\ (X\ (Suc\ i))$

shows *nat-martingale* $M\ F\ X$

proof (*intro* *nat-martingale.intro* *martingale-of-set-integral-eq*)

fix $i\ j\ A$ **assume** *asm*: $i \leq j\ A \in sets\ (F\ i)$

show *set-lebesgue-integral* $M \ A \ (X \ i) = \text{set-lebesgue-integral } M \ A \ (X \ j)$ **using** *asm*
proof (*induction* $j - i$ *arbitrary*: $i \ j$)
 case 0
 then show ?*case* **using** *asm* **by** *simp*
next
 case (*Suc* n)
 hence *: $n = j - \text{Suc } i$ **by** *linarith*
 have $\text{Suc } i \leq j$ **using** *Suc(2,3)* **by** *linarith*
 thus ?*case* **using** *sets-F-mono*[*OF - le-SucI*] *Suc(4)* *Suc(1)*[*OF **] **by** (*auto*
intro: assms(2)[*THEN trans*])
 qed
qed (*simp add: integrable*)

lemma (*in nat-sigma-finite-adapted-process*) *martingale-nat*:
 assumes *integrable*: $\bigwedge i. \text{integrable } M \ (X \ i)$
 and $\bigwedge i. AE \ \xi \text{ in } M. X \ i \ \xi = \text{cond-exp } M \ (F \ i) \ (X \ (\text{Suc } i)) \ \xi$
 shows *nat-martingale* $M \ F \ X$
proof (*unfold-locales*)
 fix $i \ j :: \text{nat}$ **assume** *asm*: $i \leq j$
 show $AE \ \xi \text{ in } M. X \ i \ \xi = \text{cond-exp } M \ (F \ i) \ (X \ j) \ \xi$ **using** *asm*
 proof (*induction* $j - i$ *arbitrary*: $i \ j$)
 case 0
 hence $j = i$ **by** *simp*
 thus ?*case* **using** *sigma-finite-subalgebra.cond-exp-F-meas*[*OF - integrable adapted*,
THEN AE-symmetric] **by** *blast*
 next
 case (*Suc* n)
 have $j: j = \text{Suc } (n + i)$ **using** *Suc* **by** *linarith*
 have $n: n = n + i - i$ **using** *Suc* **by** *linarith*
 have *: $AE \ \xi \text{ in } M. \text{cond-exp } M \ (F \ (n + i)) \ (X \ j) \ \xi = X \ (n + i) \ \xi$ **unfolding**
j **using** *assms(2)*[*THEN AE-symmetric*] **by** *blast*
 have $AE \ \xi \text{ in } M. \text{cond-exp } M \ (F \ i) \ (X \ j) \ \xi = \text{cond-exp } M \ (F \ i) \ (\text{cond-exp } M$
 $(F \ (n + i)) \ (X \ j)) \ \xi$ **by** (*intro cond-exp-nested-subalg integrable subalg, simp add:*
subalgebra-def sets-F-mono)
 hence $AE \ \xi \text{ in } M. \text{cond-exp } M \ (F \ i) \ (X \ j) \ \xi = \text{cond-exp } M \ (F \ i) \ (X \ (n + i))$
 ξ **using** *cond-exp-cong-AE*[*OF integrable-cond-exp integrable **] **by** *force*
 thus ?*case* **using** *Suc(1)*[*OF n*] **by** *fastforce*
 qed
qed (*simp add: integrable*)

lemma (*in nat-sigma-finite-adapted-process*) *martingale-of-cond-exp-diff-Suc-eq-zero*:
 assumes *integrable*: $\bigwedge i. \text{integrable } M \ (X \ i)$
 and $\bigwedge i. AE \ \xi \text{ in } M. \text{cond-exp } M \ (F \ i) \ (\lambda \xi. X \ (\text{Suc } i) \ \xi - X \ i \ \xi) \ \xi = 0$
 shows *nat-martingale* $M \ F \ X$
proof (*intro martingale-nat integrable*)
 fix i
 show $AE \ \xi \text{ in } M. X \ i \ \xi = \text{cond-exp } M \ (F \ i) \ (X \ (\text{Suc } i)) \ \xi$ **using** *cond-exp-diff*[*OF*
integrable(1,1), of i Suc i i] *cond-exp-F-meas*[*OF integrable adapted, of i*] *assms(2)*[*of*

$i]$ by *fastforce*
 qed

7.9 Discrete Time Submartingales

lemma (in *nat-submartingale*) *predictable-mono*:
 assumes *nat-predictable-process* $M F X i \leq j$
 shows $AE \xi \text{ in } M. X i \xi \leq X j \xi$
 using *assms*(2)
proof (*induction* $j - i$ arbitrary: $i j$)
 case 0
 then show ?case by *simp*
next
 case (*Suc* n)
 hence *: $n = j - Suc i$ by *linarith*
interpret S : *nat-adapted-process* $M F \lambda i. X (Suc i)$ by (*intro nat-predictable-process.adapted-Suc*
assms)
 have $Suc i \leq j$ using *Suc*(2,3) by *linarith*
 thus ?case using *Suc*(1)[*OF* *] *S.adapted*[*of i*] *submartingale-property*[*OF -*
le-SucI, *of i*] *sigma-finite-subalgebra.cond-exp-F-meas*[*OF - integrable*, *of F i Suc*
i] by *fastforce*
 qed

lemma (in *nat-sigma-finite-adapted-process-linorder*) *submartingale-of-set-integral-le-Suc*:
 assumes *integrable*: $\bigwedge i. \text{integrable } M (X i)$
 and $\bigwedge A i. A \in F i \implies \text{set-lebesgue-integral } M A (X i) \leq \text{set-lebesgue-integral}$
 $M A (X (Suc i))$
 shows *nat-submartingale* $M F X$
proof (*intro nat-submartingale.intro submartingale-of-set-integral-le*)
 fix $i j A$ assume *asm*: $i \leq j \wedge A \in \text{sets } (F i)$
 show $\text{set-lebesgue-integral } M A (X i) \leq \text{set-lebesgue-integral } M A (X j)$ using
asm
proof (*induction* $j - i$ arbitrary: $i j$)
 case 0
 then show ?case using *asm* by *simp*
next
 case (*Suc* n)
 hence *: $n = j - Suc i$ by *linarith*
 have $Suc i \leq j$ using *Suc*(2,3) by *linarith*
 thus ?case using *sets-F-mono*[*OF - le-SucI*] *Suc*(4) *Suc*(1)[*OF* *] by (*auto*
intro: assms(2)[*THEN order-trans*])
 qed
 qed (*simp add: integrable*)

lemma (in *nat-sigma-finite-adapted-process-linorder*) *submartingale-nat*:
 assumes *integrable*: $\bigwedge i. \text{integrable } M (X i)$
 and $\bigwedge i. AE \xi \text{ in } M. X i \xi \leq \text{cond-exp } M (F i) (X (Suc i)) \xi$
 shows *nat-submartingale* $M F X$
 using *subalg integrable assms*(2)

by (*intro submartingale-of-set-integral-le-Suc ord-le-eq-trans*[*OF set-integral-mono-AE-banach cond-exp-set-integral*[*symmetric*]], *simp*)
 (*meson in-mono integrable-mult-indicator set-integrable-def subalgebra-def, meson integrable-cond-exp in-mono integrable-mult-indicator set-integrable-def subalgebra-def, fast+*)

lemma (*in nat-sigma-finite-adapted-process-linorder*) *submartingale-of-cond-exp-diff-Suc-nonneg*:
assumes *integrable*: $\bigwedge i. \text{integrable } M \ (X \ i)$
and $\bigwedge i. AE \ \xi \text{ in } M. \text{cond-exp } M \ (F \ i) \ (\lambda \xi. X \ (Suc \ i) \ \xi - X \ i \ \xi) \ \xi \geq 0$
shows *nat-submartingale* $M \ F \ X$
proof (*intro submartingale-nat integrable*)
fix i
show $AE \ \xi \text{ in } M. X \ i \ \xi \leq \text{cond-exp } M \ (F \ i) \ (X \ (Suc \ i)) \ \xi$ **using** *cond-exp-diff*[*OF integrable(1,1), of i Suc i i*] *cond-exp-F-meas*[*OF integrable adapted, of i*] *assms(2)*[*of i*] **by** *fastforce*
qed

lemma (*in nat-submartingale-linorder*) *partial-sum-scaleR*:
assumes *nat-adapted-process* $M \ F \ C \ \bigwedge i. AE \ \xi \text{ in } M. 0 \leq C \ i \ \xi \ \bigwedge i. AE \ \xi \text{ in } M. C \ i \ \xi \leq R$
shows *nat-submartingale* $M \ F \ (\lambda n \ \xi. \sum i < n. C \ i \ \xi *_{\mathbb{R}} (X \ (Suc \ i) \ \xi - X \ i \ \xi))$
proof–
interpret C : *nat-adapted-process* $M \ F \ C$ **by** (*rule assms*)
interpret C' : *nat-adapted-process* $M \ F \ \lambda i \ \xi. C \ (i - 1) \ \xi *_{\mathbb{R}} (X \ i \ \xi - X \ (i - 1) \ \xi)$ **by** (*intro nat-adapted-process.intro adapted-process.scaleR-right-adapted adapted-process.diff-adapted, unfold-locales*) (*auto intro: adaptedD C.adaptedD*) +
interpret C'' : *nat-adapted-process* $M \ F \ \lambda n \ \xi. \sum i < n. C \ i \ \xi *_{\mathbb{R}} (X \ (Suc \ i) \ \xi - X \ i \ \xi)$ **by** (*rule C'.partial-sum-Suc-adapted*[*unfolded diff-Suc-1*])
interpret S : *nat-sigma-finite-adapted-process-linorder* $M \ F \ (\lambda n \ \xi. \sum i < n. C \ i \ \xi *_{\mathbb{R}} (X \ (Suc \ i) \ \xi - X \ i \ \xi)) \ \dots$
have *integrable* $M \ (\lambda x. C \ i \ x *_{\mathbb{R}} (X \ (Suc \ i) \ x - X \ i \ x))$ **for** i **using** *assms(2,3)*[*of i*] **by** (*intro Bochner-Integration.integrable-bound*[*OF integrable-scaleR-right, OF Bochner-Integration.integrable-diff, OF integrable(1,1), of R Suc i i*]) (*auto simp add: mult-mono*)
moreover have $AE \ \xi \text{ in } M. 0 \leq \text{cond-exp } M \ (F \ i) \ (\lambda \xi. (\sum i < Suc \ i. C \ i \ \xi *_{\mathbb{R}} (X \ (Suc \ i) \ \xi - X \ i \ \xi)) - (\sum i < i. C \ i \ \xi *_{\mathbb{R}} (X \ (Suc \ i) \ \xi - X \ i \ \xi))) \ \xi$ **for** i
using *sigma-finite-subalgebra.cond-exp-measurable-scaleR*[*OF - calculation - C.adapted, of i*]
cond-exp-diff-nonneg[*OF - le-SucI, OF - order.refl, of i*] *assms(2,3)*[*of i*]
by (*fastforce simp add: scaleR-nonneg-nonneg integrable*)
ultimately show *?thesis* **by** (*intro S.submartingale-of-cond-exp-diff-Suc-nonneg Bochner-Integration.integrable-sum, blast+*)
qed

lemma (*in nat-submartingale-linorder*) *partial-sum-scaleR'*:
assumes *nat-predictable-process* $M \ F \ C \ \bigwedge i. AE \ \xi \text{ in } M. 0 \leq C \ i \ \xi \ \bigwedge i. AE \ \xi \text{ in } M. C \ i \ \xi \leq R$
shows *nat-submartingale* $M \ F \ (\lambda n \ \xi. \sum i < n. C \ (Suc \ i) \ \xi *_{\mathbb{R}} (X \ (Suc \ i) \ \xi - X \ i \ \xi))$

proof –
interpret C : *nat-predictable-process* $M F C$ **by** (*rule assms*)
interpret $Suc\text{-}C$: *nat-adapted-process* $M F \lambda i. C (Suc\ i)$ **using** $C.adapted\text{-}Suc$.
show $?thesis$ **by** (*intro partial-sum-scaleR*[*of - R*] *assms*) (*intro-locales*)
qed

7.10 Discrete Time Supermartingales

lemma (*in nat-supermartingale*) *predictable-mono*:
assumes *nat-predictable-process* $M F X\ i \leq j$
shows $AE\ \xi\ in\ M. X\ i\ \xi \geq X\ j\ \xi$
using *assms*(2)
proof (*induction j - i arbitrary: i j*)
case 0
then show $?case$ **by** *simp*
next
case ($Suc\ n$)
hence $*$: $n = j - Suc\ i$ **by** *linarith*
interpret S : *nat-adapted-process* $M F \lambda i. X (Suc\ i)$ **by** (*intro nat-predictable-process.adapted-Suc assms*)
have $Suc\ i \leq j$ **using** $Suc(2,3)$ **by** *linarith*
thus $?case$ **using** $Suc(1)[OF\ *]$ $S.adapted[of\ i]$ *supermartingale-property*[$OF - le\text{-}SucI, of\ i]$ *sigma-finite-subalgebra.cond-exp-F-meas*[$OF - integrable, of\ F\ i\ Suc\ i$] **by** *fastforce*
qed

lemma (*in nat-sigma-finite-adapted-process-linorder*) *supermartingale-of-set-integral-ge-Suc*:
assumes *integrable*: $\bigwedge i. integrable\ M\ (X\ i)$
and $\bigwedge A\ i. A \in F\ i \implies set\text{-}lebesgue\text{-}integral\ M\ A\ (X\ i) \geq set\text{-}lebesgue\text{-}integral\ M\ A\ (X\ (Suc\ i))$
shows *nat-supermartingale* $M F X$
proof –
interpret $-$: *adapted-process* $M F 0 - X$ **by** (*rule uminus-adapted*)
interpret *uminus-X*: *nat-sigma-finite-adapted-process-linorder* $M F -X$..
note $*$ = *set-integral-uminus*[*unfolded set-integrable-def, OF integrable-mult-indicator*[$OF - integrable$]]
have *nat-supermartingale* $M F (-(-\ X))$
using *ord-eq-le-trans*[$OF\ * ord\text{-}le\text{-}eq\text{-}trans[OF\ le\text{-}imp\text{-}neg\text{-}le[OF\ assms(2)]\ *[symmetric]]$]
subalgebras
by (*intro nat-supermartingale.intro submartingale.uminus nat-submartingale.axioms uminus-X.submartingale-of-set-integral-le-Suc*)
(clarsimp simp add: fun-Compl-def subalgebra-def integrable | fastforce)+
thus $?thesis$ **unfolding** *fun-Compl-def* **by** *simp*
qed

lemma (*in nat-sigma-finite-adapted-process-linorder*) *supermartingale-nat*:
assumes *integrable*: $\bigwedge i. integrable\ M\ (X\ i)$
and $\bigwedge i. AE\ \xi\ in\ M. X\ i\ \xi \geq cond\text{-}exp\ M\ (F\ i)\ (X\ (Suc\ i))\ \xi$
shows *nat-supermartingale* $M F X$


```

proof –
  interpret -: adapted-process  $M\ F\ 0\ -X$  by (rule uminus-adapted)
  interpret uminus-X: nat-sigma-finite-adapted-process-linorder  $M\ F\ -X$  ..
  have  $AE\ \xi\ \text{in } M. -X\ i\ \xi \leq \text{cond-exp } M\ (F\ i)\ (\lambda x. -X\ (Suc\ i)\ x)\ \xi$  for  $i$  using
assms(2) cond-exp-uminus[OF integrable, of i Suc i] by force
  hence nat-supermartingale  $M\ F\ (-(-X))$  by (intro nat-supermartingale.intro
submartingale.uminus nat-submartingale.axioms uminus-X.submartingale-nat) (auto
simp add: fun-Compl-def integrable)
  thus ?thesis unfolding fun-Compl-def by simp
qed

lemma (in nat-sigma-finite-adapted-process-linorder) supermartingale-of-cond-exp-diff-Suc-le-zero:
  assumes integrable:  $\bigwedge i. \text{integrable } M\ (X\ i)$ 
  and  $\bigwedge i. AE\ \xi\ \text{in } M. \text{cond-exp } M\ (F\ i)\ (\lambda \xi. X\ (Suc\ i)\ \xi - X\ i\ \xi)\ \xi \leq 0$ 
  shows nat-supermartingale  $M\ F\ X$ 
proof (intro supermartingale-nat integrable)
  fix  $i$ 
  show  $AE\ \xi\ \text{in } M. X\ i\ \xi \geq \text{cond-exp } M\ (F\ i)\ (X\ (Suc\ i))\ \xi$  using cond-exp-diff[OF
integrable(1,1), of i Suc i i] cond-exp-F-meas[OF integrable adapted, of i] assms(2)[of
i] by fastforce
qed

end

```

```

theory Example-Coin-Toss
imports Martingale HOL-Probability.Stream-Space HOL-Probability.Probability-Mass-Function
begin

```

8 Example: Coin Toss

Consider a coin-tossing game, where the coin lands on heads with probability $p \in [0, 1]$. Assume that the gambler wins a fixed amount $c > 0$ on a heads outcome and loses the same amount c on a tails outcome. Let $(X_n)_{n \in \mathbb{N}}$ be a stochastic process, where X_n denotes the gamblers fortune after the n -th coin toss. Then, we have the following three cases.

1. If $p = 1/2$, it means the coin is fair and has an equal chance of landing heads or tails. In this case, the gambler, on average, neither wins nor loses money over time. The expected value of the gamblers fortune stays the same over time. Therefore, $(X_n)_{n \in \mathbb{N}}$ is a martingale.
2. If $p \geq 1/2$, it means the coin is biased in favor of heads. In this case, the gambler is more likely to win money on each bet. Over time, the gamblers fortune tends to increase on average. Therefore, $(X_n)_{n \in \mathbb{N}}$ is a submartingale.
3. If $p \leq 1/2$, it means the coin is biased in favor of tails. In this scenario, the gambler is more likely to lose money on each bet. Over time,

the gamblers fortune decreases on average. Therefore, $(X_n)_{n \in \mathbb{N}}$ is a super-martingale.

To formalize this example, we first consider a probability space consisting of infinite sequences of coin tosses.

definition *bernoulli-stream* :: *real* \Rightarrow (*bool stream*) *measure* **where**
bernoulli-stream *p* = *stream-space* (*measure-pmf* (*bernoulli-pmf* *p*))

lemma *space-bernoulli-stream[simp]*: *space* (*bernoulli-stream* *p*) = *UNIV* **by** (*simp add: bernoulli-stream-def space-stream-space*)

We define the fortune of the player at time *n* to be the number of heads minus number of tails.

definition *fortune* :: *nat* \Rightarrow *bool stream* \Rightarrow *real* **where**
fortune *n* = ($\lambda s. \sum b \leftarrow \text{stake } (\text{Suc } n) \ s. \text{ if } b \text{ then } 1 \text{ else } -1$)

definition *toss* :: *nat* \Rightarrow *bool stream* \Rightarrow *real* **where**
toss *n* = ($\lambda s. \text{ if } \text{snth } s \ n \text{ then } 1 \text{ else } -1$)

lemma *toss-indicator-def*: *toss* *n* = *indicator* {*s*. *s* !! *n*} – *indicator* {*s*. $\neg s$!! *n*}
unfolding *toss-def indicator-def* **by** *force*

lemma *range-toss*: *range* (*toss* *n*) = {–1, 1}

proof –
have *sconst True* !! *n* **by** *simp*
moreover **have** $\neg \text{sconst False}$!! *n* **by** *simp*
ultimately **have** $\exists x. x$!! *n* $\exists x. \neg x$!! *n* **by** *blast+*
thus *?thesis* **unfolding** *toss-def image-def* **by** *auto*
qed

lemma *vimage-toss*: *toss* *n* – ‘*A* = (*if* 1 \in *A* *then* {*s*. *s* !! *n*} *else* { }) \cup (*if* –1 \in *A* *then* {*s*. $\neg s$!! *n*} *else* { })

unfolding *vimage-def toss-def* **by** *auto*

lemma *fortune-Suc*: *fortune* (*Suc* *n*) *s* = *fortune* *n* *s* + *toss* (*Suc* *n*) *s*
by (*induction* *n* *arbitrary: s*) (*simp add: fortune-def toss-def*) +

lemma *fortune-toss-sum*: *fortune* *n* *s* = ($\sum i \in \{..n\}. \text{toss } i \ s$)
by (*induction* *n* *arbitrary: s*) (*simp add: fortune-def toss-def, simp add: fortune-Suc*)

lemma *fortune-bound*: *norm* (*fortune* *n* *s*) \leq *Suc* *n* **by** (*induction* *n*) (*force simp add: fortune-toss-sum toss-def*) +

Our definition of *bernoulli-stream* constitutes a probability space.

interpretation *prob-space bernoulli-stream* *p* **unfolding** *bernoulli-stream-def* **by**
(*simp add: measure-pmf.prob-space-axioms prob-space.prob-space-stream-space*)

abbreviation *toss-filtration* $p \equiv \text{nat-natural-filtration } (\text{bernoulli-stream } p) \text{ toss}$

The stochastic process *toss* is adapted to the filtration it generates.

interpretation *toss*: *nat-adapted-process bernoulli-stream p nat-natural-filtration (bernoulli-stream p) toss toss*
by (*intro nat-adapted-process.intro stochastic-process.adapted-process-natural-filtration*)
(unfold-locales, auto simp add: toss-def bernoulli-stream-def)

Similarly, the stochastic process *fortune* is adapted to the filtration generated by the tosses.

interpretation *fortune*: *nat-finite-adapted-process-linorder bernoulli-stream p nat-natural-filtration (bernoulli-stream p) toss fortune*

proof –

show *nat-finite-adapted-process-linorder (bernoulli-stream p) (toss-filtration p) fortune*

unfolding *fortune-toss-sum*

by (*intro nat-finite-adapted-process-linorder.intro*
finite-adapted-process-linorder.intro
finite-adapted-process-order.intro
finite-adapted-process.intro
toss.partial-sum-adapted[folded atMost-atLeast0]) intro-locales

qed

lemma *integrable-toss*: *integrable (bernoulli-stream p) (toss n)*

using *toss.random-variable*

by (*intro Bochner-Integration.integrable-bound[OF integrable-const[of - 1 :: real]]*)
(auto simp add: toss-def)

lemma *integrable-fortune*: *integrable (bernoulli-stream p) (fortune n) using fortune-bound*

by (*intro Bochner-Integration.integrable-bound[OF integrable-const[of - Suc n] fortune.random-variable]*) *auto*

We provide the following lemma to explicitly calculate the probability of events in this probability space.

lemma *measure-bernoulli-stream-snth-pred*:

assumes $0 \leq p$ **and** $p \leq 1$ **and** *finite J*

shows *prob p {w ∈ space (bernoulli-stream p). $\forall j \in J. P j = w !! j$ } = $p^{\wedge \text{card } (J \cap \text{Collect } P)} * (1 - p)^{\wedge (\text{card } (J - \text{Collect } P)}$*

proof –

let *?PiE = $(\Pi_E i \in J. \text{if } P i \text{ then } \{\text{True}\} \text{ else } \{\text{False}\})$*

have *product-prob-space ($\lambda \cdot \text{measure-pmf (bernoulli-pmf } p)$)* **by** *unfold-locales*

hence **: to-stream – ‘ $\{s. \forall i \in J. P i = s !! i\} = \{s. \forall i \in J. P i = s i\}$ ’* **using** *assms* **by** (*simp add: to-stream-def*)

also have *... = prod-emb UNIV ($\lambda \cdot \text{measure-pmf (bernoulli-pmf } p)$) J ?PiE*

proof –

{

```

    fix s assume (∀ i ∈ J. P i = s i)
    hence (∀ i ∈ J. P i = s i) = (s ∈ prod-emb UNIV (λ-. measure-pmf (bernoulli-pmf
p))) J ?PiE)
    by (subst prod-emb-iff[of s]) (smt (verit, best) not-def assms(3) id-def
PiE-eq-singleton UNIV-I extensional-UNIV insert-iff singletonD space-measure-pmf)
  }
  moreover
  {
    fix s assume ¬(∀ i ∈ J. P i = s i)
    then obtain i where i ∈ J P i ≠ s i by blast
    hence (∀ i ∈ J. P i = s i) = (s ∈ prod-emb UNIV (λ-. measure-pmf (bernoulli-pmf
p))) J ?PiE)
    by (simp add: restrict-def prod-emb-iff[of s]) (smt (verit, ccfv-SIG) PiE-mem
assms(3) id-def insert-iff singleton-iff)
  }
  ultimately show ?thesis by auto
qed
  finally have integ: (to-stream - ' {s. ∀ i ∈ J. P i = s !! i}) = prod-emb UNIV
(λ-. measure-pmf (bernoulli-pmf p)) J ?PiE .
  let ?M = (PiM UNIV (λ-. measure-pmf (bernoulli-pmf p)))
  have emeasure (bernoulli-stream p) {s ∈ space (bernoulli-stream p). ∀ i ∈ J. P i
= s !! i} = emeasure ?M (to-stream - ' {s. ∀ i ∈ J. P i = s !! i})
  using assms emeasure-distr[of to-stream ?M (vimage-algebra (streams (space
(measure-pmf (bernoulli-pmf p)))) (!) ?M) {s. ∀ i ∈ J. P i = s !! i}, symmetric]
measurable-to-stream[of (measure-pmf (bernoulli-pmf p))]
  by (simp only: bernoulli-stream-def stream-space-def *, simp add: space-PiM )
(smt (verit, best) emeasure-notin-sets in-vimage-algebra inf-top.right-neutral sets-distr
vimage-Collect)
  also have ... = emeasure ?M (prod-emb UNIV (λ-. measure-pmf (bernoulli-pmf
p)) J ?PiE) using integ by (simp add: space-PiM)
  also have ... = (∏ i ∈ J. emeasure (measure-pmf (bernoulli-pmf p)) (if P i then
{True} else {False}))
  by (subst emeasure-PiM-emb) (auto simp add: prob-space-measure-pmf assms(3))
  also have ... = (∏ i ∈ J ∩ Collect P. ennreal p) * (∏ i ∈ J - Collect P. ennreal
(1 - p))
  unfolding emeasure-pmf-single[of bernoulli-pmf p True, unfolded pmf-bernoulli-True[OF
assms(1,2)], symmetric]
emeasure-pmf-single[of bernoulli-pmf p False, unfolded pmf-bernoulli-False[OF
assms(1,2)], symmetric]
  by (simp add: prod.Int-Diff[OF assms(3), of - Collect P])
  also have ... = p ^ card (J ∩ Collect P) * (1 - p) ^ card (J - Collect P) using
assms by (simp add: prod-ennreal ennreal-mult' ennreal-power)
  finally show ?thesis using assms by (intro measure-eq-emeasure-eq-ennreal)
auto
qed

```

lemma

assumes $0 \leq p$ **and** $p \leq 1$

shows $\text{measure-bernoulli-stream-snth: prob } p \{w \in \text{space (bernoulli-stream } p)\}. w$

```

!! i} = p
  and measure-bernoulli-stream-neg-snth: prob p {w ∈ space (bernoulli-stream p)}.
¬w !! i} = 1 - p
  using measure-bernoulli-stream-snth-pred[OF assms, of {i} λx. True]
    measure-bernoulli-stream-snth-pred[OF assms, of {i} λx. False] by auto

```

Now we can express the expected value of a single coin toss.

```

lemma integral-toss:
  assumes 0 ≤ p p ≤ 1
  shows expectation p (toss n) = 2 * p - 1
proof -
  have [simp]: {s. s !! n} ∈ events p using measurable-snth[THEN measurable-sets,
of {True} measure-pmf (bernoulli-pmf p) n, folded bernoulli-stream-def]
  by (simp add: vimage-def)
  have expectation p (toss n) = Bochner-Integration.simple-bochner-integral (bernoulli-stream
p) (toss n)
  using toss.random-variable[of n, THEN measurable-sets]
  by (intro simple-bochner-integrable-eq-integral[symmetric] simple-bochner-integrable.intros)
(auto simp add: toss-def simple-function-def image-def)
  also have ... = p - prob p {s. ¬ s !! n} unfolding simple-bochner-integral-def
using measure-bernoulli-stream-snth[OF assms]
  by (simp add: range-toss, simp add: toss-def)
  also have ... = p - (1 - prob p {s. s !! n}) by (subst prob-compl[symmetric],
auto simp add: Collect-neg-eq Compl-eq-Diff-UNIV)
  finally show ?thesis using measure-bernoulli-stream-snth[OF assms] by simp
qed

```

Now, we show that the tosses are independent from one another.

```

lemma indep-vars-toss:
  assumes 0 ≤ p p ≤ 1
  shows indep-vars p (λ-. borel) toss {0..}
proof (subst indep-vars-def, intro conjI indep-sets-sigma)
  {
    fix A J assume asm: J ≠ {} finite J ∀j∈J. A j ∈ {toss j - ' A ∩ space
(bernoulli-stream p) | A. A ∈ borel}
    hence ∀j∈J. ∃ B ∈ borel. A j = toss j - ' B ∩ space (bernoulli-stream p) by
auto
    then obtain B where B-is: A j = toss j - ' B j ∩ space (bernoulli-stream p)
B j ∈ borel if j ∈ J for j by metis

    have prob p (∩ (A - ' J)) = (∏j∈J. prob p (A j))
    proof cases

```

We consider the case where there is a zero probability event.

```

  assume ∃j ∈ J. 1 ∉ B j ∧ -1 ∉ B j
  then obtain j where j-is: j ∈ J 1 ∉ B j -1 ∉ B j by blast
  hence A-j-empty: A j = {} using B-is by (force simp add: toss-def vimage-def)
  hence ∩ (A - ' J) = {} using j-is by blast
  moreover have prob p (A j) = 0 using A-j-empty by simp

```

ultimately show *?thesis* **using** *j-is asm(2)* **by** *auto*
next

We now assume all events have positive probability.

assume $\neg(\exists j \in J. 1 \notin B j \wedge -1 \notin B j)$
hence $*$: $1 \in B j \vee -1 \in B j$ **if** $j \in J$ **for** j **using** *that* **by** *blast*

define J' **where** [*simp*]: $J' = \{j \in J. (1 \in B j) \longleftrightarrow (-1 \notin B j)\}$
hence *toss j w* $\in B j \longleftrightarrow (1 \in B j) = w !! j$ **if** $j \in J'$ **for** $w j$ **using** *that*
unfolding *toss-def* **by** *simp*
hence $(\bigcap (A \text{ ' } J')) = \{w \in \text{space } (\text{bernoulli-stream } p). \forall j \in J'. (1 \in B j) = w !! j\}$ **using** *B-is* **by** *force*
hence *prob-J'*: $\text{prob } p (\bigcap (A \text{ ' } J')) = p \wedge \text{card } (J' \cap \{j. 1 \in B j\}) * (1 - p) \wedge \text{card } (J' - \{j. 1 \in B j\})$
using *measure-bernoulli-stream-snth-pred[OF assms finite-subset[OF - asm(2)], of J' $\lambda j. 1 \in B j$]* **by** *auto*

The index set J' consists of the indices of all non-trivial events.

have *A-j-True*: $A j = \{w \in \text{space } (\text{bernoulli-stream } p). w !! j\}$ **if** $j \in J' \cap \{j. 1 \in B j\}$ **for** j
using *that* **by** (*auto simp add: toss-def B-is(1) split: if-splits*)

have *A-j-False*: $A j = \{w \in \text{space } (\text{bernoulli-stream } p). \neg w !! j\}$ **if** $j \in J' - \{j. 1 \in B j\}$ **for** j
using *that B-is* **by** (*auto simp add: toss-def*)

have *A-j-top*: $A j = \text{space } (\text{bernoulli-stream } p)$ **if** $j \in J - J'$ **for** j **using** *that*
 $*$ **by** (*auto simp add: B-is toss-def*)

hence $\bigcap (A \text{ ' } J) = \bigcap (A \text{ ' } J')$ **by** *auto*

hence $\text{prob } p (\bigcap (A \text{ ' } J)) = \text{prob } p (\bigcap (A \text{ ' } J'))$ **by** *presburger*

also have $\dots = (\prod_{j \in J' \cap \{j. 1 \in B j\}} \text{prob } p (A j)) * (\prod_{j \in J' - \{j. 1 \in B j\}} \text{prob } p (A j))$

by (*simp only: prob-J' A-j-True A-j-False measure-bernoulli-stream-snth[OF assms] measure-bernoulli-stream-neg-snth[OF assms] cong: prod.cong*) *simp*

also have $\dots = (\prod_{j \in J'} \text{prob } p (A j))$ **using** *asm(2)* **by** (*intro prod.Int-Diff[symmetric]*) *auto*

also have $\dots = (\prod_{j \in J'} \text{prob } p (A j)) * (\prod_{j \in J - J'} \text{prob } p (A j))$ **using** *A-j-top prob-space* **by** *simp*

also have $\dots = (\prod_{j \in J} \text{prob } p (A j))$ **using** *asm(2)* **by** (*metis (no-types, lifting) J'-def mem-Collect-eq mult.commute prod.subset-diff subsetI*)

finally show *?thesis* .

qed

}

thus *indep-sets p* ($\lambda i. \{ \text{toss } i - ' A \cap \text{space } (\text{bernoulli-stream } p) \mid A. A \in \text{sets borel} \} \{0..\}$) **using** *measurable-sets[OF toss.random-variable]*

by (*intro indep-setsI subsetI*) *fastforce*

qed (*simp, intro Int-stableI, simp, metis sets.Int vimage-Int*)

The fortune of a player is a martingale (resp. sub- or supermartingale) with

respect to the filtration generated by the coin tosses.

theorem *fortune-martingale*:

assumes $p = 1/2$
shows *nat-martingale* (*bernoulli-stream* p) (*toss-filtration* p) *fortune*
using *cond-exp-indep*[*OF* *fortune.subalg* *indep-set-natural-filtration* *integrable-toss*,
OF *zero-order*(1) *lessI* *indep-vars-toss*, of p]
integral-toss *assms*
by (*intro* *fortune.martingale-of-cond-exp-diff-Suc-eq-zero* *integrable-fortune*)
(force simp add: fortune-toss-sum)

theorem *fortune-submartingale*:

assumes $1/2 \leq p \leq 1$
shows *nat-submartingale* (*bernoulli-stream* p) (*toss-filtration* p) *fortune*
proof (*intro* *fortune.submartingale-of-cond-exp-diff-Suc-nonneg* *integrable-fortune*)
fix n
show *AE* ξ in *bernoulli-stream* p . $0 \leq \text{cond-exp}$ (*bernoulli-stream* p) (*toss-filtration*
 p n) ($\lambda \xi. \text{fortune} (\text{Suc } n) \xi - \text{fortune } n \xi$) ξ
using *cond-exp-indep*[*OF* *fortune.subalg* *indep-set-natural-filtration* *integrable-toss*,
OF *zero-order*(1) *lessI* *indep-vars-toss*, of p n]
integral-toss[of p *Suc* n] *assms*
by (*force simp add: fortune-toss-sum*)
qed

theorem *fortune-supermartingale*:

assumes $0 \leq p \leq 1/2$
shows *nat-supermartingale* (*bernoulli-stream* p) (*toss-filtration* p) *fortune*
proof (*intro* *fortune.supermartingale-of-cond-exp-diff-Suc-le-zero* *integrable-fortune*)
fix n
show *AE* ξ in *bernoulli-stream* p . $0 \geq \text{cond-exp}$ (*bernoulli-stream* p) (*toss-filtration*
 p n) ($\lambda \xi. \text{fortune} (\text{Suc } n) \xi - \text{fortune } n \xi$) ξ
using *cond-exp-indep*[*OF* *fortune.subalg* *indep-set-natural-filtration* *integrable-toss*,
OF *zero-order*(1) *lessI* *indep-vars-toss*, of p n]
integral-toss[of p *Suc* n] *assms*
by (*force simp add: fortune-toss-sum*)
qed

end

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