

Low Degree Hypergraphs

Ata Keskin

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Abstract

Martingale poopies: [1], I love my bebis alara!

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1 Introduction

Alara is best bebis ever!

```
theory Measure-Space-Addendum
imports HOL-Analysis.Measure-Space
begin
```

1.1 Sigma Algebra Generated by a Family of Functions

definition *sigma-gen* :: 'a set \Rightarrow 'b measure \Rightarrow ('a \Rightarrow 'b) set \Rightarrow 'a measure **where**
sigma-gen Ω N $S \equiv \text{sigma } \Omega (\bigcup f \in S. \{f - 'A \cap \Omega \mid A. A \in N\})$

lemma [*simp*]:

shows *sets-sigma-gen*: *sets* (*sigma-gen* Ω N S) = *sigma-sets* $\Omega (\bigcup f \in S. \{f - 'A \cap \Omega \mid A. A \in N\})$
and *space-sigma-gen*: *space* (*sigma-gen* Ω N S) = Ω
by (*auto simp add: sigma-gen-def sets-measure-of-conv space-measure-of-conv*)

lemma *measurable-sigma-gen*:

assumes $f \in S$ $f \in \Omega \rightarrow \text{space } N$
shows $f \in \text{sigma-gen } \Omega$ N $S \rightarrow_M N$
using *assms* **by** (*intro measurableI, auto*)

lemma *measurable-sigma-gen-singleton*:

assumes $f \in \Omega \rightarrow \text{space } N$
shows $f \in \text{sigma-gen } \Omega$ N $\{f\} \rightarrow_M N$
using *assms* *measurable-sigma-gen* **by** *blast*

lemma *measurable-iff-contains-sigma-gen*:

shows $(f \in M \rightarrow_M N) \longleftrightarrow f \in \text{space } M \rightarrow \text{space } N \wedge \text{sigma-gen } (\text{space } M) N$
 $\{f\} \subseteq M$
proof (*standard, goal-cases*)
case 1
hence $f \in \text{space } M \rightarrow \text{space } N$ **using** *measurable-space* **by** *fast*
thus ?case **unfolding** *sets-sigma-gen* **by** (*simp, intro sigma-algebra.sigma-sets-subset,*
(*blast intro: sets.sigma-algebra-axioms measurable-sets[OF 1]*)+)
next
case 2
thus ?case **using** *measurable-mono*[*OF - refl - space-sigma-gen, of N M*] *measurable-sigma-gen-singleton* **by** *fast*
qed

lemma *measurable-iff-contains-sigma-gen'*:

shows $(S \subseteq M \rightarrow_M N) \longleftrightarrow S \subseteq \text{space } M \rightarrow \text{space } N \wedge \text{sigma-gen } (\text{space } M) N S \subseteq M$
proof (*standard, goal-cases*)
case 1
hence *subset*: $S \subseteq \text{space } M \rightarrow \text{space } N$ **using** *measurable-space* **by** *fast*
have $\{f - 'A \cap \text{space } M \mid A. A \in N\} \subseteq M$ **if** $f \in S$ **for** f **using** *measur-*

```

able-iff-contains-sigma-gen[unfolded sets-sigma-gen, of f] 1 subset that by blast
  then show ?case unfolding sets-sigma-gen using sets.sigma-algebra-axioms by
(simp add: subset, intro sigma-algebra.sigma-sets-subset, blast+)
next
  case 2
  hence subset:  $S \subseteq \text{space } M \rightarrow \text{space } N$  by simp
  show ?case
  proof (standard, goal-cases)
    case (1 x)
    have sigma-gen (space M) N {x}  $\subseteq M$  by (metis (no-types, lifting) 1 2
sets-sigma-gen SUP-le-iff sigma-sets-le-sets-iff singletonD)
    thus ?case using measurable-iff-contains-sigma-gen subset[THEN subsetD, OF
1] by fast
  qed
qed

end
theory Elementary-Metric-Spaces-Addendum
imports HOL-Analysis.Elementary-Metric-Spaces HOL-Analysis.Bochner-Integration
begin

lemma diameter-comp-strict-mono:
  fixes s :: nat  $\Rightarrow$  'a :: real-normed-vector
  assumes strict-mono r bounded {s i | i. r n  $\leq$  i}
  shows diameter {s (r i) | i. n  $\leq$  i}  $\leq$  diameter {s i | i. r n  $\leq$  i}
proof (rule diameter-subset)
  show {s (r i) | i. n  $\leq$  i}  $\subseteq$  {s i | i. r n  $\leq$  i} using assms(1) monotoneD
strict-mono-mono by fastforce
qed (intro assms(2))

lemma diameter-bounded-bound':
  fixes S :: 'a :: metric-space set
  assumes S: bdd-above (case-prod dist ' (S  $\times$  S)) x  $\in$  S y  $\in$  S
  shows dist x y  $\leq$  diameter S
proof -
  have (x,y)  $\in$  S  $\times$  S using S by auto
  then have dist x y  $\leq$  (SUP (x,y)  $\in$  S  $\times$  S. dist x y) by (rule cSUP-upper2[OF
assms(1)]) simp
  with  $\langle x \in S \rangle$  show ?thesis by (auto simp: diameter-def)
qed

lemma bounded-imp-dist-bounded:
  assumes bounded (range s)
  shows bounded (( $\lambda(i, j). \text{dist } (s \ i) (s \ j)$ ) ' ({n..}  $\times$  {n..}))
  using bounded-dist-comp[OF bounded-fst bounded-snd, OF bounded-Times(1,1)[OF
assms(1,1)]] by (rule bounded-subset, force)

lemma cauchy-iff-diameter-tends-to-zero-and-bounded:
  fixes s :: nat  $\Rightarrow$  'a :: real-normed-vector

```

shows $\text{Cauchy } s \iff ((\lambda n. \text{diameter } \{s\ i \mid i. i \geq n\}) \longrightarrow 0 \wedge \text{bounded (range } s))$
proof –
have $\{s\ i \mid i. N \leq i\} \neq \{\}$ **for** N **by** *blast*
hence *diameter-SUP*: $\text{diameter } \{s\ i \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s\ i) (s\ j))$ **for** N **unfolding** *diameter-def* **by** (*auto intro!*: *arg-cong[of - - Sup]*)
show *?thesis*
proof ((*standard ; clarsimp*), *goal-cases*)
case 1
have $\exists N. \forall n \geq N. \text{norm } (\text{diameter } \{s\ i \mid i. n \leq i\}) < e$ **if** *e-pos*: $e > 0$ **for** e
proof –
obtain N **where** *dist-less*: $\text{dist } (s\ n) (s\ m) < (1/2) * e$ **if** $n \geq N \ m \geq N$ **for** $n\ m$ **using** 1 *CauchyD e-pos dist-norm* **by** (*metis mult-pos-pos zero-less-divide-iff zero-less-numeral zero-less-one*)
{
fix r **assume** $r \geq N$
hence $\text{dist } (s\ n) (s\ m) < (1/2) * e$ **if** $n \geq r \ m \geq r$ **for** $n\ m$ **using** *dist-less*
that by *simp*
hence $(\text{SUP } (i, j) \in \{r..\} \times \{r..\}. \text{dist } (s\ i) (s\ j)) \leq (1/2) * e$ **by** (*intro cSup-least*) *fastforce* +
also have $\dots < e$ **using** *e-pos* **by** *simp*
finally have $\text{diameter } \{s\ i \mid i. r \leq i\} < e$ **using** *diameter-SUP* **by** *presburger*
}
moreover have $\text{diameter } \{s\ i \mid i. r \leq i\} \geq 0$ **for** r **unfolding** *diameter-SUP*
using *bounded-imp-dist-bounded[OF cauchy-imp-bounded, THEN bounded-imp-bdd-above, OF 1]* **by** (*intro cSup-upper2, auto*)
ultimately show *?thesis* **by** *auto*
qed
thus *?case* **using** *cauchy-imp-bounded[OF 1]* **by** (*simp add: LIMSEQ-iff*)
next
case 2
have $\exists N. \forall n \geq N. \forall m \geq N. \text{dist } (s\ n) (s\ m) < e$ **if** *e-pos*: $e > 0$ **for** e
proof –
obtain N **where** *diam-less*: $\text{diameter } \{s\ i \mid i. r \leq i\} < e$ **if** $r \geq N$ **for** r
using *LIMSEQ-D 2(1) e-pos* **by** *fastforce*
{
fix $n\ m$ **assume** $n \geq N \ m \geq N$
hence $\text{dist } (s\ n) (s\ m) < e$ **using** *cSUP-lessD[OF bounded-imp-dist-bounded[THEN bounded-imp-bdd-above], OF 2(2) diam-less[unfolded diameter-SUP]]* **by** *auto*
}
thus *?thesis* **by** *blast*
qed
then show *?case* **by** (*intro CauchyI, simp add: dist-norm*)
qed
qed
context
fixes $s\ r :: \text{nat} \Rightarrow 'a \Rightarrow 'b :: \{\text{second-countable-topology, real-normed-vector, banach}\}$ **and** M

assumes *bounded*: $\bigwedge x. x \in \text{space } M \implies \text{bounded } (\text{range } (\lambda i. s \ i \ x))$
begin

lemma *borel-measurable-diameter*:

assumes [*measurable*]: $\bigwedge i. (s \ i) \in \text{borel-measurable } M$

shows $(\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\}) \in \text{borel-measurable } M$

proof –

have $\{s \ i \ x \mid i. N \leq i\} \neq \{\}$ **for** $x \ N$ **by** *blast*

hence *diameter-SUP*: $\text{diameter } \{s \ i \ x \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s \ i \ x) (s \ j \ x))$ **for** $x \ N$ **unfolding** *diameter-def* **by** (*auto intro!*: *arg-cong[of - - Sup]*)

have *case-prod dist* ‘ $(\{s \ i \ x \mid i. n \leq i\} \times \{s \ i \ x \mid i. n \leq i\}) = ((\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) ' (\{n..\} \times \{n..\}))$ ’ **for** x **by** *fast*

hence *: $(\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\}) = (\lambda x. \text{Sup } ((\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) ' (\{n..\} \times \{n..\})))$ **using** *diameter-SUP* **by** (*simp add: case-prod-beta'*)

have *bounded* $((\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) ' (\{n..\} \times \{n..\}))$ **if** $x \in \text{space } M$ **for** x **by** (*rule bounded-imp-dist-bounded[OF bounded, OF that]*)

hence *bdd*: *bdd-above* $((\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) ' (\{n..\} \times \{n..\}))$ **if** $x \in \text{space } M$ **for** x **using** *that bounded-imp-bdd-above* **by** *presburger*

have *fst* $p \in \text{borel-measurable } M$ *snd* $p \in \text{borel-measurable } M$ **if** $p \in s ' \{n..\} \times s ' \{n..\}$ **for** p **using** *that* **by** *fastforce+*

hence $(\lambda x. \text{fst } p \ x - \text{snd } p \ x) \in \text{borel-measurable } M$ **if** $p \in s ' \{n..\} \times s ' \{n..\}$ **for** p **using** *that borel-measurable-diff* **by** *simp*

hence $(\lambda x. \text{case } p \text{ of } (f, g) \Rightarrow \text{dist } (f \ x) (g \ x)) \in \text{borel-measurable } M$ **if** $p \in s ' \{n..\} \times s ' \{n..\}$ **for** p **unfolding** *dist-norm* **using** *that* **by** *measurable*

moreover **have** *countable* $(s ' \{n..\} \times s ' \{n..\})$ **by** (*intro countable-SIGMA countable-image, auto*)

ultimately show *?thesis* **unfolding** * **by** (*auto intro!*: *borel-measurable-cSUP bdd*)

qed

lemma *integrable-bound-diameter*:

fixes $f :: 'a \Rightarrow \text{real}$

assumes *integrable* $M \ f$

and [*measurable*]: $\bigwedge i. (s \ i) \in \text{borel-measurable } M$

and $\bigwedge x \ i. x \in \text{space } M \implies \text{norm } (s \ i \ x) \leq f \ x$

shows *integrable* $M \ (\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\})$

proof –

have $\{s \ i \ x \mid i. N \leq i\} \neq \{\}$ **for** $x \ N$ **by** *blast*

hence *diameter-SUP*: $\text{diameter } \{s \ i \ x \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s \ i \ x) (s \ j \ x))$ **for** $x \ N$ **unfolding** *diameter-def* **by** (*auto intro!*: *arg-cong[of - - Sup]*)

{

fix x **assume** $x: x \in \text{space } M$

let $?S = (\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) ' (\{n..\} \times \{n..\})$

have *case-prod dist* ‘ $(\{s \ i \ x \mid i. n \leq i\} \times \{s \ i \ x \mid i. n \leq i\}) = (\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) ' (\{n..\} \times \{n..\})$ ’ **by** *fast*

```

    hence *: diameter {s i x | i. n ≤ i} = Sup ?S using diameter-SUP by (simp
add: case-prod-beta')

    have bounded ?S by (rule bounded-imp-dist-bounded[OF bounded[OF x]])
    hence Sup-S-nonneg: 0 ≤ Sup ?S by (auto intro!: cSup-upper2 x bounded-imp-bdd-above)

    have dist (s i x) (s j x) ≤ 2 * f x for i j by (intro dist-triangle2[THEN
order-trans, of - 0]) (metis norm-conv-dist assms(3) x add-mono mult-2)
    hence ∀ c ∈ ?S. c ≤ 2 * f x by force
    hence Sup ?S ≤ 2 * f x by (intro cSup-least, auto)
    hence norm (Sup ?S) ≤ 2 * norm (f x) using Sup-S-nonneg by auto
    also have ... = norm (2 *R f x) by simp
    finally have norm (diameter {s i x | i. n ≤ i}) ≤ norm (2 *R f x) unfolding
* .
}
    hence AE x in M. norm (diameter {s i x | i. n ≤ i}) ≤ norm (2 *R f x) by blast
    thus integrable M (λx. diameter {s i x | i. n ≤ i}) using borel-measurable-diameter
by (intro Bochner-Integration.integrable-bound[OF assms(1)[THEN integrable-scaleR-right[of
2]]], measurable)
qed
end

end
theory Bochner-Integration-Addendum
imports HOL-Analysis.Bochner-Integration
begin

```

1.2 Simple Functions

```

lemma integrable-implies-simple-function-sequence:
  fixes f :: 'a ⇒ 'b::{banach, second-countable-topology}
  assumes integrable M f
  obtains s where ∧ i. simple-function M (s i)
    and ∧ i. emeasure M {y ∈ space M. s i y ≠ 0} ≠ ∞
    and ∧ x. x ∈ space M ⇒ (λ i. s i x) ⟶ f x
    and ∧ x i. x ∈ space M ⇒ norm (s i x) ≤ 2 * norm (f x)
proof-
  have f: f ∈ borel-measurable M (∫+x. norm (f x) ∂M) < ∞ using assms
unfolding integrable-iff-bounded by auto
  obtain s where s: ∧ i. simple-function M (s i) ∧ x. x ∈ space M ⇒ (λ i. s
i x) ⟶ f x ∧ i x. x ∈ space M ⇒ norm (s i x) ≤ 2 * norm (f x) using
borel-measurable-implies-sequence-metric[OF f(1)] unfolding norm-conv-dist by
metis
  {
    fix i
    have (∫+x. norm (s i x) ∂M) ≤ (∫+x. ennreal (2 * norm (f x)) ∂M) using
s by (intro nn-integral-mono, auto)
    also have ... < ∞ using f by (simp add: nn-integral-cmult ennreal-mult-less-top
ennreal-mult)
  }

```

finally have *sbi: Bochner-Integration.simple-bochner-integrable* M (*s i*) **using**
s **by** (*intro simple-bochner-integrableI-bounded*) *auto*
hence *emeasure* $M \{y \in \text{space } M. s \ i \ y \neq 0\} \neq \infty$ **by** (*auto intro: inte-*
grableI-simple-bochner-integrable simple-bochner-integrable.cases)
}
thus *?thesis* **using** *that s* **by** *blast*
qed

lemma *banach-simple-function-indicator-representation*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes f : *simple-function* $M \ f$ **and** x : $x \in \text{space } M$
shows $f \ x = (\sum y \in f \text{ ` } \text{space } M. \text{indicator } (f \text{ - ` } \{y\} \cap \text{space } M) \ x \ *_R \ y)$
(is ?l = ?r)
proof –
have $?r = (\sum y \in f \text{ ` } \text{space } M.$
(if $y = f \ x$ *then* $\text{indicator } (f \text{ - ` } \{y\} \cap \text{space } M) \ x \ *_R \ y$ *else* $0))$ **by** (*auto intro!:*
sum.cong)
also have $\dots = \text{indicator } (f \text{ - ` } \{f \ x\} \cap \text{space } M) \ x \ *_R \ f \ x$ **using** *assms* **by** (*auto*
dest: simple-functionD)
also have $\dots = f \ x$ **using** x **by** (*auto simp: indicator-def*)
finally show *?thesis* **by** *auto*
qed

lemma *banach-simple-function-indicator-representation-AE*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes f : *simple-function* $M \ f$
shows *AE* x *in* $M. f \ x = (\sum y \in f \text{ ` } \text{space } M. \text{indicator } (f \text{ - ` } \{y\} \cap \text{space } M) \ x$
 $*_R \ y)$
by (*metis (mono-tags, lifting) AE-I2 banach-simple-function-indicator-representation*
f)

lemmas *simple-function-scaleR*[*intro*] = *simple-function-compose2*[**where** $h = (*_R)$]

lemma *integrable-simple-function*:
assumes *simple-function* $M \ f$ *emeasure* $M \ \{y \in \text{space } M. f \ y \neq 0\} \neq \infty$
shows *integrable* $M \ f$
using *assms* *has-bochner-integral-simple-bochner-integrable integrable.simps sim-*
ple-bochner-integrable.simps **by** *blast*

lemma *simple-integrable-function-induct*[*consumes 2, case-names cong indicator*
add, induct set: simple-function]:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes f : *simple-function* $M \ f$ *emeasure* $M \ \{y \in \text{space } M. f \ y \neq 0\} \neq \infty$
assumes *cong*: $\bigwedge f \ g. \text{simple-function } M \ f \implies \text{emeasure } M \ \{y \in \text{space } M. f \ y \neq 0\} \neq \infty \implies \text{simple-function } M \ g \implies \text{emeasure } M \ \{y \in \text{space } M. g \ y \neq 0\} \neq \infty$
 $\implies (\bigwedge x. x \in \text{space } M \implies f \ x = g \ x) \implies P \ f \implies P \ g$
assumes *indicator*: $\bigwedge A \ y. A \in \text{sets } M \implies \text{emeasure } M \ A < \infty \implies P \ (\lambda x.$
indicator $A \ x \ *_R \ y)$
assumes *add*: $\bigwedge f \ g. \text{simple-function } M \ f \implies \text{emeasure } M \ \{y \in \text{space } M. f \ y \neq$

$0\} \neq \infty \implies$
 $\infty \implies$
 $(g\ z)) \implies$
 $P\ f \implies P\ g \implies P\ (\lambda x. f\ x + g\ x)$
shows $P\ f$
proof –
let $?f = \lambda x. (\sum y \in f.\ 'space\ M. \text{indicat-real}\ (f - '\{y\} \cap space\ M)\ x *_R y)$
have $f\text{-ae-eq}: f\ x = ?f\ x$ **if** $x \in space\ M$ **for** x **using** $\text{banach-simple-function-indicator-representation}[OF\ f(1)\ \text{that}]$.
moreover have $\text{emeasure}\ M\ \{y \in space\ M. ?f\ y \neq 0\} \neq \infty$ **by** $(metis\ (no-types,\ lifting)\ \text{Collect-cong}\ \text{calculation}\ f(2))$
moreover have $P\ (\lambda x. \sum y \in S. \text{indicat-real}\ (f - '\{y\} \cap space\ M)\ x *_R y)$
 $\text{simple-function}\ M\ (\lambda x. \sum y \in S. \text{indicat-real}\ (f - '\{y\} \cap space\ M)\ x$
 $*_R y)$
 $\text{emeasure}\ M\ \{y \in space\ M. (\sum x \in S. \text{indicat-real}\ (f - '\{x\} \cap space\ M)\ y *_R x) \neq 0\} \neq \infty$
if $S \subseteq f.\ 'space\ M$ **for** S **using** $\text{simple-functionD}(1)[OF\ \text{assms}(1),$
 $THEN\ rev\text{-finite-subset},\ OF\ \text{that}]$ **that**
proof $(induction\ rule:\ \text{finite-induct})$
case $empty$
 $\{$
case 1
then show $?case$ **using** $\text{indicator}[of\ \{\}]$ **by force**
next
case 2
then show $?case$ **by force**
next
case 3
then show $?case$ **by force**
 $\}$
next
case $(insert\ x\ F)$
have $(f - '\{x\} \cap space\ M) \subseteq \{y \in space\ M. f\ y \neq 0\}$ **if** $x \neq 0$ **using** $that$ **by**
 $blast$
moreover have $\{y \in space\ M. f\ y \neq 0\} = space\ M - (f - '\{0\} \cap space\ M)$
by $blast$
moreover have $space\ M - (f - '\{0\} \cap space\ M) \in sets\ M$ **using** $\text{simple-functionD}(2)[OF\ f(1)]$ **by** $blast$
ultimately have $\text{fin-0}: \text{emeasure}\ M\ (f - '\{x\} \cap space\ M) < \infty$ **if** $x \neq 0$
using $that$ **by** $(metis\ \text{emeasure-mono}\ f(2)\ \text{infinity-ennreal-def}\ top.\text{not-eq-extremum}\ top.\text{unique})$
hence $\text{fin-1}: \text{emeasure}\ M\ \{y \in space\ M. \text{indicat-real}\ (f - '\{x\} \cap space\ M)\ y *_R$
 $x \neq 0\} \neq \infty$ **if** $x \neq 0$ **by** $(metis\ (mono-tags,\ lifting)\ \text{emeasure-mono}\ f(1)\ \text{indica-}$
 $\text{tor-simps}(2)\ \text{linorder-not-less}\ mem\text{-Collect-eq}\ \text{scaleR-eq-0-iff}\ \text{simple-functionD}(2)$
 $\text{subsetI}\ \text{that})$
have $*$: $(\sum y \in insert\ x\ F. \text{indicat-real}\ (f - '\{y\} \cap space\ M)\ x *_R y) = (\sum y \in F.$


```

indicat-real (f - ' {y} ∩ space M) xa *R y) + indicat-real (f - ' {x} ∩ space M)
xa *R x for xa by (metis (no-types, lifting) Diff-empty Diff-insert0 add commute
insert.hyps(1) insert.hyps(2) sum.insert-remove)
  have **: {y ∈ space M. (∑ x ∈ insert x F. indicat-real (f - ' {x} ∩ space M) y
*_R x) ≠ 0} ⊆ {y ∈ space M. (∑ x ∈ F. indicat-real (f - ' {x} ∩ space M) y *_R x)
≠ 0} ∪ {y ∈ space M. indicat-real (f - ' {x} ∩ space M) y *_R x ≠ 0} unfolding
* by fastforce
  {
    case 1
    hence x: x ∈ f ' space M and F: F ⊆ f ' space M by auto
    show ?case
    proof (cases x = 0)
      case True
      then show ?thesis unfolding * using insert(3)[OF F] by simp
    next
      case False
      have norm-argument: norm ((∑ y ∈ F. indicat-real (f - ' {y} ∩ space M) z
*_R y) + indicat-real (f - ' {x} ∩ space M) z *_R x) = norm (∑ y ∈ F. indicat-real
(f - ' {y} ∩ space M) z *_R y) + norm (indicat-real (f - ' {x} ∩ space M) z *_R x)
if z: z ∈ space M for z
      proof (cases f z = x)
        case True
        have indicat-real (f - ' {y} ∩ space M) z *_R y = 0 if y ∈ F for y using
True insert(2) z that 1 unfolding indicator-def by force
        hence (∑ y ∈ F. indicat-real (f - ' {y} ∩ space M) z *_R y) = 0 by (meson
sum.neutral)
        then show ?thesis by force
      next
        case False
        then show ?thesis by force
      qed
      show ?thesis using False simple-functionD(2)[OF f(1)] insert(3,5)[OF F]
simple-function-scaleR fin-0 fin-1 by (subst *, subst add, subst simple-function-sum)
      (blast intro: norm-argument indicator)+
      qed
    next
      case 2
      hence x: x ∈ f ' space M and F: F ⊆ f ' space M by auto
      show ?case
      proof (cases x = 0)
        case True
        then show ?thesis unfolding * using insert(4)[OF F] by simp
      next
        case False
        then show ?thesis unfolding * using insert(4)[OF F] simple-functionD(2)[OF
f(1)] by fast
      qed
    next
      case 3

```

hence $x: x \in f \text{ ' space } M$ **and** $F: F \subseteq f \text{ ' space } M$ **by** *auto*
show *?case*
proof (*cases* $x = 0$)
 case *True*
 then show *?thesis unfolding * using insert(5)[OF F]* **by** *simp*
next
 case *False*
 have $\text{emeasure } M \{y \in \text{space } M. (\sum x \in \text{insert } x F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \leq \text{emeasure } M (\{y \in \text{space } M. (\sum x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\})$
 using *** simple-functionD(2)[OF insert(4)[OF F]] simple-functionD(2)[OF f(1)]* **by** (*intro emeasure-mono, force+*)
 also have $\dots \leq \text{emeasure } M \{y \in \text{space } M. (\sum x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} + \text{emeasure } M \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\}$
 using *simple-functionD(2)[OF insert(4)[OF F]] simple-functionD(2)[OF f(1)]* **by** (*intro emeasure-subadditive, force+*)
 also have $\dots < \infty$ **using** *insert(5)[OF F] fin-1[OF False]* **by** (*simp add: less-top*)
 finally show *?thesis* **by** *simp*
 qed
}
qed
moreover have *simple-function* $M (\lambda x. \sum y \in f \text{ ' space } M. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y)$ **using** *calculation* **by** *blast*
moreover have $P (\lambda x. \sum y \in f \text{ ' space } M. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y)$ **using** *calculation* **by** *blast*
ultimately show *?thesis* **by** (*intro cong[OF - - f(1,2)], blast, presburger+*)
qed

lemma *simple-integrable-function-induct-nonneg[consumes 3, case-names cong indicator add, induct set: simple-function]:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes f : *simple-function* $M f$ $\text{emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty \wedge x. x \in \text{space } M \longrightarrow f x \geq 0$

assumes *cong*: $\wedge f g. \text{simple-function } M f \Longrightarrow \text{emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty \Longrightarrow (\wedge x. x \in \text{space } M \Longrightarrow f x \geq 0) \Longrightarrow \text{simple-function } M g \Longrightarrow \text{emeasure } M \{y \in \text{space } M. g y \neq 0\} \neq \infty \Longrightarrow (\wedge x. x \in \text{space } M \Longrightarrow g x \geq 0) \Longrightarrow (\wedge x. x \in \text{space } M \Longrightarrow f x = g x) \Longrightarrow P f \Longrightarrow P g$

assumes *indicator*: $\wedge A y. y \geq 0 \Longrightarrow A \in \text{sets } M \Longrightarrow \text{emeasure } M A < \infty \Longrightarrow P (\lambda x. \text{indicator } A x *_R y)$

assumes *add*: $\wedge f g. (\wedge x. x \in \text{space } M \Longrightarrow f x \geq 0) \Longrightarrow \text{simple-function } M f \Longrightarrow \text{emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty \Longrightarrow$

$(\wedge x. x \in \text{space } M \Longrightarrow g x \geq 0) \Longrightarrow \text{simple-function } M g \Longrightarrow \text{emeasure } M \{y \in \text{space } M. g y \neq 0\} \neq \infty \Longrightarrow$

$(\wedge z. z \in \text{space } M \Longrightarrow \text{norm } (f z + g z) = \text{norm } (f z) + \text{norm } (g z)) \Longrightarrow$

$P f \implies P g \implies P (\lambda x. f x + g x)$

shows $P f$

proof –

let $?f = \lambda x. (\sum y \in f \text{ ' space } M. \text{ indicat-real } (f - \text{' } \{y\} \cap \text{space } M) x *_R y)$

have $f\text{-ae-eq}: f x = ?f x$ **if** $x \in \text{space } M$ **for** x **using** *banach-simple-function-indicator-representation*[*OF f(1) that*] .

moreover have $\text{emeasure } M \{y \in \text{space } M. ?f y \neq 0\} \neq \infty$ **by** (*metis (no-types, lifting) Collect-cong calculation f(2)*)

moreover have $P (\lambda x. \sum y \in S. \text{ indicat-real } (f - \text{' } \{y\} \cap \text{space } M) x *_R y)$

simple-function $M (\lambda x. \sum y \in S. \text{ indicat-real } (f - \text{' } \{y\} \cap \text{space } M) x *_R y)$

$\text{emeasure } M \{y \in \text{space } M. (\sum x \in S. \text{ indicat-real } (f - \text{' } \{x\} \cap \text{space } M) y *_R x) \neq 0\} \neq \infty$

$\bigwedge x. x \in \text{space } M \implies 0 \leq (\sum y \in S. \text{ indicat-real } (f - \text{' } \{y\} \cap \text{space } M) x *_R y)$

if $S \subseteq f \text{ ' space } M$ **for** S **using** *simple-functionD(1)*[*OF assms(1), THEN rev-finite-subset, OF that*] **that**

proof (*induction rule: finite-subset-induct'*)

case *empty*

{

case 1

then show $?case$ **using** *indicator*[*of 0 {}*] **by force**

next

case 2

then show $?case$ **by force**

next

case 3

then show $?case$ **by force**

next

case 4

then show $?case$ **by force**

}

next

case (*insert x F*)

have $(f - \text{' } \{x\} \cap \text{space } M) \subseteq \{y \in \text{space } M. f y \neq 0\}$ **if** $x \neq 0$ **using** *that* **by** *blast*

moreover have $\{y \in \text{space } M. f y \neq 0\} = \text{space } M - (f - \text{' } \{0\} \cap \text{space } M)$

by *blast*

moreover have $\text{space } M - (f - \text{' } \{0\} \cap \text{space } M) \in \text{sets } M$ **using** *simple-functionD(2)*[*OF f(1)*] **by** *blast*

ultimately have $\text{fin-0}: \text{emeasure } M (f - \text{' } \{x\} \cap \text{space } M) < \infty$ **if** $x \neq 0$ **using** *that* **by** (*metis emeasure-mono f(2) infinity-ennreal-def top.not-eq-extremum top-unique*)

hence $\text{fin-1}: \text{emeasure } M \{y \in \text{space } M. \text{ indicat-real } (f - \text{' } \{x\} \cap \text{space } M) y *_R x \neq 0\} \neq \infty$ **if** $x \neq 0$ **by** (*metis (mono-tags, lifting) emeasure-mono f(1) indicator-simps(2) linorder-not-less mem-Collect-eq scaleR-eq-0-iff simple-functionD(2) subsetI that*)

have $\text{nonneg-}x: x \geq 0$ **using** *insert f* **by** *blast*

have *: $(\sum y \in \text{insert } x \ F. \text{indicat-real } (f - \{y\} \cap \text{space } M) \ x \ast_R y) =$
 $(\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) \ x \ast_R y) + \text{indicat-real } (f - \{x\} \cap$
 $\text{space } M) \ x \ast_R x$ **for** x **by** (*metis* (*no-types*, *lifting*) *add.commute insert.hyps(1)*
insert.hyps(4) sum.insert)
have **: $\{y \in \text{space } M. (\sum x \in \text{insert } x \ F. \text{indicat-real } (f - \{x\} \cap \text{space } M) \ y$
 $\ast_R x) \neq 0\} \subseteq \{y \in \text{space } M. (\sum x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) \ y \ast_R x)$
 $\neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) \ y \ast_R x \neq 0\}$ **unfolding**
***** **by** *fastforce*
{
 case 1
 show ?*case*
 proof (*cases* $x = 0$)
 case *True*
 then show ?*thesis* **unfolding** ***** **using** *insert* **by** *simp*
 next
 case *False*
 have *norm-argument*: $\text{norm } ((\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) \ z$
 $\ast_R y) + \text{indicat-real } (f - \{x\} \cap \text{space } M) \ z \ast_R x) = \text{norm } (\sum y \in F. \text{indicat-real}$
 $(f - \{y\} \cap \text{space } M) \ z \ast_R y) + \text{norm } (\text{indicat-real } (f - \{x\} \cap \text{space } M) \ z \ast_R x)$
if $z \in \text{space } M$ **for** z
 proof (*cases* $f \ z = x$)
 case *True*
 have $\text{indicat-real } (f - \{y\} \cap \text{space } M) \ z \ast_R y = 0$ **if** $y \in F$ **for** y **using**
True insert z that 1 **unfolding** *indicator-def* **by** *force*
 hence $(\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) \ z \ast_R y) = 0$ **by** (*meson*
sum.neutral)
 thus ?*thesis* **by** *force*
 qed (*force*)
 show ?*thesis* **using** *False fin-0 fin-1 f norm-argument* **by** (*subst* ***, *subst* *add*,
presburger add: insert, intro insert, intro insert, fastforce simp add: indicator-def
intro!: insert(2) f(3), auto intro!: indicator insert nonneg-x)
 qed
next
 case 2
 show ?*case*
 proof (*cases* $x = 0$)
 case *True*
 then show ?*thesis* **unfolding** ***** **using** *insert* **by** *simp*
 next
 case *False*
 then show ?*thesis* **unfolding** ***** **using** *insert simple-functionD(2)[OF f(1)]*
by *fast*
 qed
next
 case 3
 show ?*case*
 proof (*cases* $x = 0$)
 case *True*
 then show ?*thesis* **unfolding** ***** **using** *insert* **by** *simp*

next
case *False*
have $\text{emeasure } M \{y \in \text{space } M. (\sum x \in \text{insert } x \ F. \text{indicat-real } (f - ' \{x\} \cap \text{space } M) \ y *_{\mathbb{R}} x) \neq 0\} \leq \text{emeasure } M (\{y \in \text{space } M. (\sum x \in F. \text{indicat-real } (f - ' \{x\} \cap \text{space } M) \ y *_{\mathbb{R}} x) \neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - ' \{x\} \cap \text{space } M) \ y *_{\mathbb{R}} x \neq 0\})$
using *** simple-functionD(2)[OF insert(6)] simple-functionD(2)[OF f(1)] insert.IH(2)* **by** (*intro emeasure-mono, blast, simp*)
also have $\dots \leq \text{emeasure } M \{y \in \text{space } M. (\sum x \in F. \text{indicat-real } (f - ' \{x\} \cap \text{space } M) \ y *_{\mathbb{R}} x) \neq 0\} + \text{emeasure } M \{y \in \text{space } M. \text{indicat-real } (f - ' \{x\} \cap \text{space } M) \ y *_{\mathbb{R}} x \neq 0\}$
using *simple-functionD(2)[OF insert(6)] simple-functionD(2)[OF f(1)]*
by (*intro emeasure-subadditive, force+*)
also have $\dots < \infty$ **using** *insert(7) fin-1[OF False]* **by** (*simp add: less-top*)
finally show *?thesis* **by** *simp*
qed
next
case ($\frac{1}{4} \xi$)
thus *?case using insert nonneg-x f(3)* **by** (*auto simp add: scaleR-nonneg-nonneg intro: sum-nonneg*)
}
qed
moreover have *simple-function* $M (\lambda x. \sum y \in f^{-1} \text{space } M. \text{indicat-real } (f - ' \{y\} \cap \text{space } M) \ x *_{\mathbb{R}} y)$ **using** *calculation* **by** *blast*
moreover have $P (\lambda x. \sum y \in f^{-1} \text{space } M. \text{indicat-real } (f - ' \{y\} \cap \text{space } M) \ x *_{\mathbb{R}} y)$ **using** *calculation* **by** *blast*
moreover have $\bigwedge x. x \in \text{space } M \implies 0 \leq f \ x$ **using** *f(3)* **by** *simp*
ultimately show *?thesis* **by** (*intro cong[OF - - - f(1,2)], blast, blast, fast*)
presburger+
qed

proposition *integrable-induct'[consumes 1, case-names base add lim, induct pred: integrable]:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{banach, second-countable-topology}\}$
assumes *integrable* $M \ f$
assumes *base:* $\bigwedge A \ c. A \in \text{sets } M \implies \text{emeasure } M \ A < \infty \implies P (\lambda x. \text{indicator } A \ x *_{\mathbb{R}} c)$
assumes *add:* $\bigwedge f \ g. \text{integrable } M \ f \implies P \ f \implies \text{integrable } M \ g \implies P \ g \implies P (\lambda x. f \ x + g \ x)$
assumes *lim:* $\bigwedge f \ s. \text{integrable } M \ f$
 $\implies (\bigwedge i. \text{integrable } M \ (s \ i))$
 $\implies (\bigwedge i. \text{simple-function } M \ (s \ i))$
 $\implies (\bigwedge i. \text{emeasure } M \ \{y \in \text{space } M. s \ i \ y \neq 0\} \neq \infty)$
 $\implies (\bigwedge x. x \in \text{space } M \implies (\lambda i. s \ i \ x) \longrightarrow f \ x)$
 $\implies (\bigwedge i \ x. x \in \text{space } M \implies \text{norm } (s \ i \ x) \leq 2 * \text{norm } (f \ x))$
 $\implies (\bigwedge i. P \ (s \ i)) \implies P \ f$

shows $P \ f$

proof –

have $f: f \in \text{borel-measurable } M \ (\int^+ x. \text{norm } (f \ x) \ \partial M) < \infty$ **using** *assms(1)*

unfolding *integrable-iff-bounded* **by** *auto*

obtain *s* **where** *s*: $\bigwedge i. \text{simple-function } M (s\ i) \bigwedge x. x \in \text{space } M \implies (\lambda i. s\ i\ x) \longrightarrow f\ x \bigwedge i x. x \in \text{space } M \implies \text{norm } (s\ i\ x) \leq 2 * \text{norm } (f\ x)$ **using** *borel-measurable-implies-sequence-metric*[*OF* *f*(1)] **unfolding** *norm-conv-dist* **by** *metis*

```
{
  fix f A
  have [simp]:  $P (\lambda x. 0)$  using base[of {} undefined] by simp
  have  $(\bigwedge i::'b. i \in A \implies \text{integrable } M (f\ i::'a \Rightarrow 'b)) \implies (\bigwedge i. i \in A \implies P (f\ i)) \implies P (\lambda x. \sum i \in A. f\ i\ x)$  by (induct A rule: infinite-finite-induct) (auto intro!: add)
}
```

note *sum* = *this*

define *s'* **where** [*abs-def*]: $s'\ i\ z = \text{indicator } (\text{space } M)\ z\ *_R\ s\ i\ z$ **for** *i* *z*
hence *s'-eq-s*: $\bigwedge i x. x \in \text{space } M \implies s'\ i\ x = s\ i\ x$ **by** *simp*

have *sf*[*measurable*]: $\bigwedge i. \text{simple-function } M (s'\ i)$ **unfolding** *s'-def* **using** *s*(1) **by** (*intro simple-function-compose2*[**where** *h*=(**_R*)] *simple-function-indicator*) *auto*

```
{
  fix i
  have  $\bigwedge z. \{y. s'\ i\ z = y \wedge y \in s'\ i\ ' \text{space } M \wedge y \neq 0 \wedge z \in \text{space } M\} = (\text{if } z \in \text{space } M \wedge s'\ i\ z \neq 0 \text{ then } \{s'\ i\ z\} \text{ else } \{\})$  by (auto simp add: s'-def split: split-indicator)
  then have  $\bigwedge z. s'\ i = (\lambda z. \sum y \in s'\ i\ ' \text{space } M - \{0\}. \text{indicator } \{x \in \text{space } M. s'\ i\ x = y\}\ z\ *_R\ y)$  using sf by (auto simp: fun-eq-iff simple-function-def s'-def)
}
```

note *s'-eq* = *this*

show *P f*

proof (*rule* *lim*)

fix *i*

have $(\int^+ x. \text{norm } (s'\ i\ x)\ \partial M) \leq (\int^+ x. \text{ennreal } (2 * \text{norm } (f\ x))\ \partial M)$ **using** *s* **by** (*intro nn-integral-mono*) (*auto simp*: *s'-eq-s*)

also have $\dots < \infty$ **using** *f* **by** (*simp add*: *nn-integral-cmult ennreal-mult-less-top ennreal-mult*)

finally have *sbi*: *Bochner-Integration.simple-bochner-integrable* *M* (*s' i*) **using** *sf* **by** (*intro simple-bochner-integrableI-bounded*) *auto*

thus *integrable* *M* (*s' i*) *simple-function* *M* (*s' i*) *emeasure* *M* $\{y \in \text{space } M. s'\ i\ y \neq 0\} \neq \infty$ **by** (*auto intro*: *integrableI-simple-bochner-integrable simple-bochner-integrable.cases*)

```
{
  fix x assume  $x \in \text{space } M\ s'\ i\ x \neq 0$ 
  then have  $\text{emeasure } M \{y \in \text{space } M. s'\ i\ y = s'\ i\ x\} \leq \text{emeasure } M \{y \in \text{space } M. s'\ i\ y \neq 0\}$  by (intro emeasure-mono) auto
  also have  $\dots < \infty$  using sbi by (auto elim: simple-bochner-integrable.cases)
}
```

```

simp: less-top)
  finally have emeasure M {y ∈ space M. s' i y = s' i x} ≠ ∞ by simp
}
then show P (s' i) by (subst s'-eq) (auto intro!: sum base simp: less-top)

fix x assume x ∈ space M
thus (λi. s' i x) ⟶ f x using s by (simp add: s'-eq-s)
show norm (s' i x) ≤ 2 * norm (f x) using ⟨x ∈ space M⟩ s by (simp add:
s'-eq-s)
qed fact
qed

lemma finite-nn-integral-imp-ae-finite:
  fixes f :: 'a ⇒ ennreal
  assumes f ∈ borel-measurable M (∫+x. f x ∂M) < ∞
  shows AE x in M. f x < ∞
proof (rule ccontr, goal-cases)
  case 1
  let ?A = space M ∩ {x. f x = ∞}
  have *: emeasure M ?A > 0 using 1 assms(1) by (metis (mono-tags, lifting)
assms(2) eventually-mono infinity-ennreal-def nn-integral-noteq-infinite top.not-eq-extremum)
  have (∫+x. f x * indicator ?A x ∂M) = (∫+x. ∞ * indicator ?A x ∂M) by
(metis (mono-tags, lifting) indicator-inter-arith indicator-simps(2) mem-Collect-eq
mult-eq-0-iff)
  also have ... = ∞ * emeasure M ?A using assms(1) by (intro nn-integral-cmult-indicator,
simp)
  also have ... = ∞ using * by fastforce
  finally have (∫+x. f x * indicator ?A x ∂M) = ∞ .
  moreover have (∫+x. f x * indicator ?A x ∂M) ≤ (∫+x. f x ∂M) by (intro
nn-integral-mono, simp add: indicator-def)
  ultimately show ?case using assms(2) by simp
qed

lemma cauchy-L1-AE-cauchy-subseq:
  fixes s :: nat ⇒ 'a ⇒ 'b::{banach, second-countable-topology}
  assumes [measurable]: ⋀n. integrable M (s n)
  and ⋀e. e > 0 ⟹ ∃N. ∀i ≥ N. ∀j ≥ N. LINT x|M. dist (s i x) (s j x) < e
  obtains r where strict-mono r AE x in M. Cauchy (λi. s (r i) x)
proof-
  have ∃r. ∀n. (∀i ≥ r n. ∀j ≥ r n. LINT x|M. dist (s i x) (s j x) < (1 / 2) ^ n)
  ∧ (r (Suc n) > r n)
  proof (intro dependent-nat-choice, goal-cases)
    case 1
    then show ?case using assms(2) by presburger
  next
    case (2 x n)
    obtain N where *: LINT x|M. dist (s i x) (s j x) < (1 / 2) ^ Suc n if i ≥ N
    j ≥ N for i j using assms(2)[of (1 / 2) ^ Suc n] by auto
    {

```

fix $i\ j$ **assume** $i \geq \max N\ (Suc\ x)\ j \geq \max N\ (Suc\ x)$
hence $integral^L\ M\ (\lambda x. dist\ (s\ i\ x)\ (s\ j\ x)) < (1 / 2) \wedge Suc\ n$ **using** $*$ **by**
fastforce
}
then show $?case$ **by** *fastforce*
qed
then obtain r **where** *strict-mono: strict-mono* r **and** $\forall i \geq r\ n. \forall j \geq r\ n. LINT\ x | M. dist\ (s\ i\ x)\ (s\ j\ x) < (1 / 2) \wedge n$ **for** n **using** *strict-mono-Suc-iff* **by** *blast*
hence $r-is: LINT\ x | M. dist\ (s\ (r\ (Suc\ n))\ x)\ (s\ (r\ n)\ x) < (1 / 2) \wedge n$ **for** n
by (*simp add: strict-mono-leD*)

define g **where** $g = (\lambda n\ x. (\sum i \leq n. ennreal\ (dist\ (s\ (r\ (Suc\ i))\ x)\ (s\ (r\ i)\ x))))$
define g' **where** $g' = (\lambda x. \sum i. ennreal\ (dist\ (s\ (r\ (Suc\ i))\ x)\ (s\ (r\ i)\ x)))$

have *integrable-g*: $(\int^+ x. g\ n\ x\ \partial M) < 2$ **for** n
proof –
have $(\int^+ x. g\ n\ x\ \partial M) = (\int^+ x. (\sum i \leq n. ennreal\ (dist\ (s\ (r\ (Suc\ i))\ x)\ (s\ (r\ i)\ x))))\ \partial M)$ **using** $g-def$ **by** *simp*
also have $\dots = (\sum i \leq n. (\int^+ x. ennreal\ (dist\ (s\ (r\ (Suc\ i))\ x)\ (s\ (r\ i)\ x))\ \partial M))$ **by** (*intro nn-integral-sum, simp*)
also have $\dots = (\sum i \leq n. LINT\ x | M. dist\ (s\ (r\ (Suc\ i))\ x)\ (s\ (r\ i)\ x))$ **unfolding** *dist-norm* **using** *assms(1)* **by** (*subst nn-integral-eq-integral[OF integrable-norm], auto*)
also have $\dots < ennreal\ (\sum i \leq n. (1 / 2) \wedge i)$ **by** (*intro ennreal-lessI[OF sum-pos sum-strict-mono[OF finite-atMost - r-is]], auto*)
also have $\dots \leq ennreal\ 2$ **unfolding** *sum-gp0[of 1 / 2 n]* **by** (*intro ennreal-leI, auto*)
finally show $(\int^+ x. g\ n\ x\ \partial M) < 2$ **by** *simp*
qed

have *integrable-g'*: $(\int^+ x. g'\ x\ \partial M) \leq 2$
proof –
have *incseq* $(\lambda n. g\ n\ x)$ **for** x **by** (*intro incseq-SucI, auto simp add: g-def ennreal-leI*)
hence *convergent* $(\lambda n. g\ n\ x)$ **for** x **unfolding** *convergent-def* **using** *LIM-SEQ-incseq-SUP* **by** *fast*
hence $(\lambda n. g\ n\ x) \longrightarrow g'\ x$ **for** x **unfolding** $g-def\ g'-def$ **by** (*intro summable-iff-convergent'[THEN iffD2, THEN summable-LIMSEQ]', blast*)
hence $(\int^+ x. g'\ x\ \partial M) = (\int^+ x. liminf\ (\lambda n. g\ n\ x)\ \partial M)$ **by** (*metis lim-imp-Liminf trivial-limit-sequentially*)
also have $\dots \leq liminf\ (\lambda n. \int^+ x. g\ n\ x\ \partial M)$ **by** (*intro nn-integral-liminf, simp add: g-def*)
also have $\dots \leq liminf\ (\lambda n. 2)$ **using** *integrable-g* **by** (*intro Liminf-mono*) (*simp add: order-le-less*)
also have $\dots = 2$ **using** *sequentially-bot tendsto-iff-Liminf-eq-Limsup* **by** *blast*
finally show $?thesis$.
qed
hence $AE\ x\ in\ M. g'\ x < \infty$ **by** (*intro finite-nn-integral-imp-ae-finite*) (*auto simp add: order-le-less-trans g'-def*)

moreover have summable $(\lambda i. \text{dist } (s \ (r \ (Suc \ i)) \ x) \ (s \ (r \ i) \ x))$ if $g' \ x \neq \infty$ for x using that unfolding g' -def by (intro summable-suminf-not-top, intro zero-le-dist, fastforce)

ultimately have ae-summable: AE x in M . summable $(\lambda i. s \ (r \ (Suc \ i)) \ x - s \ (r \ i) \ x)$ using summable-norm-cancel unfolding dist-norm by force

```
{
  fix x assume summable  $(\lambda i. s \ (r \ (Suc \ i)) \ x - s \ (r \ i) \ x)$ 
  hence  $(\lambda n. \sum i < n. s \ (r \ (Suc \ i)) \ x - s \ (r \ i) \ x) \longrightarrow (\sum i. s \ (r \ (Suc \ i)) \ x - s \ (r \ i) \ x)$  using summable-LIMSEQ by blast
  moreover have  $(\lambda n. (\sum i < n. s \ (r \ (Suc \ i)) \ x - s \ (r \ i) \ x)) = (\lambda n. s \ (r \ n) \ x - s \ (r \ 0) \ x)$  using sum-lessThan-telescope by fastforce
  ultimately have  $(\lambda n. s \ (r \ n) \ x - s \ (r \ 0) \ x) \longrightarrow (\sum i. s \ (r \ (Suc \ i)) \ x - s \ (r \ i) \ x)$  by argo
  hence  $(\lambda n. s \ (r \ n) \ x - s \ (r \ 0) \ x + s \ (r \ 0) \ x) \longrightarrow (\sum i. s \ (r \ (Suc \ i)) \ x - s \ (r \ i) \ x) + s \ (r \ 0) \ x$  by (intro isCont-tendsto-compose[of -  $\lambda z. z + s \ (r \ 0) \ x$ ], auto)
  hence Cauchy  $(\lambda n. s \ (r \ n) \ x)$  by (simp add: LIMSEQ-imp-Cauchy)
}
```

hence AE x in M . Cauchy $(\lambda i. s \ (r \ i) \ x)$ using ae-summable by fast
thus ?thesis by (rule that[OF strict-mono(1)])

qed

lemma integrable-max[simp, intro]:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology}\}$

assumes $fg[\text{measurable}]: \text{integrable } M \ f \ \text{integrable } M \ g$

shows $\text{integrable } M \ (\lambda x. \max (f \ x) \ (g \ x))$

proof (rule Bochner-Integration.integrable-bound)

```
{
  fix x y :: 'b
  have  $\text{norm } (\max x \ y) \leq \max (\text{norm } x) (\text{norm } y)$  by linarith
  also have  $\dots \leq \text{norm } x + \text{norm } y$  by simp
  finally have  $\text{norm } (\max x \ y) \leq \text{norm } (\text{norm } x + \text{norm } y)$  by simp
}
```

thus AE x in M . $\text{norm } (\max (f \ x) \ (g \ x)) \leq \text{norm } (\text{norm } (f \ x) + \text{norm } (g \ x))$ by simp

qed (auto intro: Bochner-Integration.integrable-add[OF integrable-norm[OF $fg(1)$] integrable-norm[OF $fg(2)$]])

lemma integrable-min[simp, intro]:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology}\}$

assumes $[measurable]: \text{integrable } M \ f \ \text{integrable } M \ g$

shows $\text{integrable } M \ (\lambda x. \min (f \ x) \ (g \ x))$

proof -

have $\text{norm } (\min (f \ x) \ (g \ x)) \leq \text{norm } (f \ x) \vee \text{norm } (\min (f \ x) \ (g \ x)) \leq \text{norm } (g \ x)$ for x by linarith

thus ?thesis by (intro integrable-bound[OF integrable-max[OF integrable-norm(1,1), OF assms]], auto)

qed

lemma *integral-nonneg-AE-banach*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes $[measurable]: f \in \text{borel-measurable } M$ **and** $\text{nonneg}: AE\ x\ \text{in } M. 0 \leq f\ x$

shows $0 \leq \text{integral}^L\ M\ f$

proof *cases*

assume *integrable*: $\text{integrable } M\ f$

hence $\text{max}: (\lambda x. \text{max } 0\ (f\ x)) \in \text{borel-measurable } M \wedge x. 0 \leq \text{max } 0\ (f\ x)$

integrable $M\ (\lambda x. \text{max } 0\ (f\ x))$ **by** *auto*

hence $0 \leq \text{integral}^L\ M\ (\lambda x. \text{max } 0\ (f\ x))$

proof *–*

obtain s **where** $*$: $\wedge i. \text{simple-function } M\ (s\ i)$

$\wedge i. \text{emeasure } M\ \{y \in \text{space } M. s\ i\ y \neq 0\} \neq \infty$

$\wedge x. x \in \text{space } M \implies (\lambda i. s\ i\ x) \longrightarrow f\ x$

$\wedge x\ i. x \in \text{space } M \implies \text{norm } (s\ i\ x) \leq 2 * \text{norm } (f\ x)$ **using**

integrable-implies-simple-function-sequence[*OF integrable*] **by** *blast*

have *simple*: $\wedge i. \text{simple-function } M\ (\lambda x. \text{max } 0\ (s\ i\ x))$ **using** $*$ **by** *fast*

have $\wedge i. \{y \in \text{space } M. \text{max } 0\ (s\ i\ y) \neq 0\} \subseteq \{y \in \text{space } M. s\ i\ y \neq 0\}$

unfolding *max-def* **by** *force*

moreover **have** $\wedge i. \{y \in \text{space } M. s\ i\ y \neq 0\} \in \text{sets } M$ **using** $*$ **by** *measurable*

ultimately **have** $\wedge i. \text{emeasure } M\ \{y \in \text{space } M. \text{max } 0\ (s\ i\ y) \neq 0\} \leq$

$\text{emeasure } M\ \{y \in \text{space } M. s\ i\ y \neq 0\}$ **by** (*rule emeasure-mono*)

hence $**:\wedge i. \text{emeasure } M\ \{y \in \text{space } M. \text{max } 0\ (s\ i\ y) \neq 0\} \neq \infty$ **using** $*(2)$

by (*auto intro: order.strict-trans1 simp add: top.not-eq-extremum*)

have $\wedge x. x \in \text{space } M \implies (\lambda i. \text{max } 0\ (s\ i\ x)) \longrightarrow \text{max } 0\ (f\ x)$ **using** $*(3)$

tendsto-max **by** *blast*

moreover **have** $\wedge x\ i. x \in \text{space } M \implies \text{norm } (\text{max } 0\ (s\ i\ x)) \leq \text{norm } (2 *_R$

$f\ x)$ **using** $*(4)$ **unfolding** *max-def* **by** *auto*

ultimately **have** *tendsto*: $(\lambda i. \text{integral}^L\ M\ (\lambda x. \text{max } 0\ (s\ i\ x))) \longrightarrow \text{integral}^L\ M\ (\lambda x. \text{max } 0\ (f\ x))$

using *borel-measurable-simple-function simple integrable* **by** (*intro integral-dominated-convergence*[*OF max(1) - integrable-norm*][*OF integrable-scaleR-right*], *of - 2 f*], *blast+*)

{

fix $h :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assume *simple-function* $M\ h$ $\text{emeasure } M\ \{y \in \text{space } M. h\ y \neq 0\} \neq \infty \wedge x.$

$x \in \text{space } M \longrightarrow h\ x \geq 0$

hence $*$: $\text{integral}^L\ M\ h \geq 0$

proof (*induct rule: simple-integrable-function-induct-nonneg*)

case (*cong f g*)

then show *?case* **using** *Bochner-Integration.integral-cong* **by** *force*

next

case (*indicator A y*)

hence $A \neq \{\} \implies y \geq 0$ **using** *sets.sets-into-space* **by** *fastforce*

then show *?case* **using** *indicator* **by** (*cases A = {}*], *auto simp add:*

```

scaleR-nonneg-nonneg)
  next
  case (add f g)
  then show ?case by (simp add: integrable-simple-function)
qed
}
thus ?thesis using LIMSEQ-le-const[OF tendsto, of 0] ** simple by fastforce
qed
also have ... = integralL M f using nonneg by (auto intro: integral-cong-AE)
finally show ?thesis .
qed (simp add: not-integrable-integral-eq)

lemma integral-mono-AE-banach:
  fixes f g :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
    dered-real-vector}
  assumes integrable M f integrable M g AE x in M. f x ≤ g x
  shows integralL M f ≤ integralL M g
  using integral-nonneg-AE-banach[of λx. g x - f x] assms Bochner-Integration.integral-diff[OF
    assms(1,2)] by force

lemma integral-mono-banach:
  fixes f g :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
    dered-real-vector}
  assumes integrable M f integrable M g ∧ x. x ∈ space M ⇒ f x ≤ g x
  shows integralL M f ≤ integralL M g
  using integral-mono-AE-banach assms by blast

end
theory Set-Integral-Addendum
  imports HOL-Analysis.Set-Integral Bochner-Integration-Addendum
  begin

lemma set-integral-scaleR-left:
  assumes A ∈ sets M c ≠ 0 ⇒ integrable M f
  shows LINT t:A|M. f t *R c = (LINT t:A|M. f t) *R c
  unfolding set-lebesgue-integral-def
  using integrable-mult-indicator[OF assms]
  by (subst integral-scaleR-left[symmetric], auto)

lemma nn-set-integral-eq-set-integral:
  assumes [measurable]:integrable M f
  and AE x ∈ A in M. 0 ≤ f x A ∈ sets M
  shows (∫+ x ∈ A. f x ∂M) = (∫ x ∈ A. f x ∂M)
proof-
  have (∫+ x. indicator A x *R f x ∂M) = (∫ x ∈ A. f x ∂M)
  unfolding set-lebesgue-integral-def using assms(2) by (intro nn-integral-eq-integral[of
    - λx. indicat-real A x *R f x], blast intro: assms integrable-mult-indicator, fastforce)
  moreover have (∫+ x. indicator A x *R f x ∂M) = (∫+ x ∈ A. f x ∂M) by (metis
    ennreal-0 indicator-simps(1) indicator-simps(2) mult.commute mult-1 mult-zero-left

```

real-scaleR-def)
ultimately show *?thesis* **by** *argo*
qed

lemma *set-integral-restrict-space*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{banach}, \text{second-countable-topology}\}$
assumes $\Omega \cap \text{space } M \in \text{sets } M$
shows $\text{set-lebesgue-integral } (\text{restrict-space } M \ \Omega) \ A \ f = \text{set-lebesgue-integral } M \ A$
 $(\lambda x. \text{indicator } \Omega \ x \ *_{\mathbb{R}} \ f \ x)$
unfolding *set-lebesgue-integral-def*
by (*subst integral-restrict-space, auto intro!; integrable-mult-indicator assms simp;*
mult.commute)

lemma *set-integral-const*:
fixes $c :: 'b :: \{\text{banach}, \text{second-countable-topology}\}$
assumes $A \in \text{sets } M \ \text{emeasure } M \ A \neq \infty$
shows $\text{set-lebesgue-integral } M \ A \ (\lambda \cdot. c) = \text{measure } M \ A \ *_{\mathbb{R}} \ c$
unfolding *set-lebesgue-integral-def*
using *assms* **by** (*metis has-bochner-integral-indicator has-bochner-integral-integral-eq*
infinity-enreal-def less-top)

lemma *set-integral-mono-banach*:
fixes $f \ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}, \text{linorder-topology}, \text{ordered-real-vector}\}$
assumes $\text{set-integrable } M \ A \ f \ \text{set-integrable } M \ A \ g$
 $\bigwedge x. x \in A \implies f \ x \leq g \ x$
shows $(\text{LINT } x:A | M. f \ x) \leq (\text{LINT } x:A | M. g \ x)$
using *assms* **unfolding** *set-integrable-def set-lebesgue-integral-def*
by (*auto intro: integral-mono-banach split: split-indicator*)

lemma *set-integral-mono-AE-banach*:
fixes $f \ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}, \text{linorder-topology}, \text{ordered-real-vector}\}$
assumes $\text{set-integrable } M \ A \ f \ \text{set-integrable } M \ A \ g \ \text{AE } x \in A \text{ in } M. f \ x \leq g \ x$
shows $\text{set-lebesgue-integral } M \ A \ f \leq \text{set-lebesgue-integral } M \ A \ g$ **using** *assms*
unfolding *set-lebesgue-integral-def* **by** (*auto simp add: set-integrable-def intro!;*
*integral-mono-AE-banach[of M $\lambda x. \text{indicator } A \ x \ *_{\mathbb{R}} \ f \ x$ $\lambda x. \text{indicator } A \ x \ *_{\mathbb{R}} \ g \ x$],*
simp add: indicator-def)

end
theory *Sigma-Finite-Measure-Addendum*
imports *Set-Integral-Addendum*
begin

lemma *balls-countable-basis*:
obtains $D :: 'a :: \{\text{metric-space}, \text{second-countable-topology}\} \text{ set}$
where *topological-basis* (*case-prod ball ' (D \times ($\mathbb{Q} \cap \{0 < ..\}$))*)
and *countable* D

```

    and  $D \neq \{\}$ 
  proof -
    obtain  $D :: 'a \text{ set}$  where  $\text{dense-subset: countable } D \ D \neq \{\}$   $[[\text{open } U; U \neq \{\}]]$ 
     $\implies \exists y \in D. y \in U$  for  $U$  using  $\text{countable-dense-exists}$  by  $\text{blast}$ 
    have  $\text{topological-basis}$   $(\text{case-prod ball } ' (D \times (\mathbb{Q} \cap \{0 < ..\})))$ 
    proof  $(\text{intro topological-basis-iff}[THEN \text{iffD2}], \text{fast}, \text{clarify})$ 
    fix  $U$  and  $x :: 'a$  assume  $\text{asm: open } U \ x \in U$ 
    obtain  $e$  where  $e: e > 0$   $\text{ball } x \ e \subseteq U$  using  $\text{asm openE}$  by  $\text{blast}$ 
    obtain  $y$  where  $y: y \in D \ y \in \text{ball } x \ (e / 3)$  using  $\text{dense-subset}(3)[OF \text{ open-ball},$ 
 $\text{of } x \ e / 3]$   $\text{centre-in-ball}[THEN \text{iffD2}, OF \text{ divide-pos-pos}[OF \text{ e(1)}, \text{ of } 3]]$  by  $\text{force}$ 
    obtain  $r$  where  $r: r \in \mathbb{Q} \cap \{e/3 < .. < e/2\}$  unfolding  $\text{Rats-def}$  using  $\text{of-rat-dense}[OF$ 
 $\text{divide-strict-left-mono}[OF - \text{e(1)}, \text{ of } 2 \ 3]$  by  $\text{auto}$ 

    have  $*$ :  $x \in \text{ball } y \ r$  using  $r \ y$  by  $(\text{simp add: dist-commute})$ 
    hence  $\text{ball } y \ r \subseteq U$  using  $r$  by  $(\text{intro order-trans}[OF - \text{e(2)}], \text{simp}, \text{metric})$ 
    moreover have  $\text{ball } y \ r \in (\text{case-prod ball } ' (D \times (\mathbb{Q} \cap \{0 < ..\})))$  using  $y(1)$ 
 $r$  by  $\text{force}$ 
    ultimately show  $\exists B' \in (\text{case-prod ball } ' (D \times (\mathbb{Q} \cap \{0 < ..\}))). x \in B' \wedge B' \subseteq$ 
 $U$  using  $*$  by  $\text{meson}$ 
    qed
    thus  $?thesis$  using  $\text{that dense-subset}$  by  $\text{blast}$ 
  qed

context  $\text{sigma-finite-measure}$ 
begin

lemma  $\text{sigma-finite-measure-induct}[case-names \text{ finite-measure}, \text{ consumes } 0]$ :
  assumes  $\bigwedge(N :: 'a \text{ measure}) \ \Omega. \text{finite-measure } N$ 
     $\implies N = \text{restrict-space } M \ \Omega$ 
     $\implies \Omega \in \text{sets } M$ 
     $\implies \text{emeasure } N \ \Omega \neq \infty$ 
     $\implies \text{emeasure } N \ \Omega \neq 0$ 
     $\implies \text{almost-everywhere } N \ Q$ 
  and  $[\text{measurable}]: \text{Measurable.pred } M \ Q$ 
shows  $\text{almost-everywhere } M \ Q$ 
proof -
    have  $*$ :  $\text{almost-everywhere } N \ Q$  if  $\text{finite-measure } N \ N = \text{restrict-space } M \ \Omega \ \Omega$ 
 $\in \text{sets } M \ \text{emeasure } N \ \Omega \neq \infty$  for  $N \ \Omega$  using  $\text{that}$  by  $(\text{cases } \text{emeasure } N \ \Omega = 0,$ 
 $\text{auto intro: emeasure-0-AE assms}(1))$ 

    obtain  $A :: \text{nat} \Rightarrow 'a \text{ set}$  where  $A: \text{range } A \subseteq \text{sets } M \ (\bigcup i. A \ i) = \text{space } M$  and
 $\text{emeasure-finite: emeasure } M \ (A \ i) \neq \infty$  for  $i$  using  $\text{sigma-finite}$  by  $\text{metis}$ 
    note  $A(1)[\text{measurable}]$ 
    have  $\text{space-restr: space } (\text{restrict-space } M \ (A \ i)) = A \ i$  for  $i$  unfolding  $\text{space-restrict-space}$ 
by  $\text{simp}$ 
    {
      fix  $i$ 
      have  $*$ :  $\{x \in A \ i \cap \text{space } M. \ Q \ x\} = \{x \in \text{space } M. \ Q \ x\} \cap (A \ i)$  by  $\text{fast}$ 
      have  $\text{Measurable.pred } (\text{restrict-space } M \ (A \ i)) \ Q$  using  $A$  by  $(\text{intro measurableI},$ 

```

```

auto simp add: space-restr intro!: sets-restrict-space-iff[THEN iffD2], measurable,
auto)
}
note this[measurable]
{
  fix i
  have finite-measure (restrict-space M (A i)) using emeasure-finite by (intro
finite-measureI, subst space-restr, subst emeasure-restrict-space, auto)
  hence emeasure (restrict-space M (A i))  $\{x \in A \ i. \neg Q \ x\} = 0$  using emea-
sure-finite by (intro AE-iff-measurable[THEN iffD1, OF - - *], measurable, subst
space-restr[symmetric], intro sets.top, auto simp add: emeasure-restrict-space)
  hence emeasure M  $\{x \in A \ i. \neg Q \ x\} = 0$  by (subst emeasure-restrict-space[symmetric],
auto)
}
hence emeasure M  $(\bigcup i. \{x \in A \ i. \neg Q \ x\}) = 0$  by (intro emeasure-UN-eq-0,
auto)
moreover have  $(\bigcup i. \{x \in A \ i. \neg Q \ x\}) = \{x \in \text{space } M. \neg Q \ x\}$  using A by
auto
ultimately show ?thesis by (intro AE-iff-measurable[THEN iffD2], auto)
qed

```

lemma averaging-theorem:

```

fixes f:: $\Rightarrow$  'b::{second-countable-topology, banach}
assumes [measurable]:integrable M f
and closed: closed S
and  $\bigwedge A. A \in \text{sets } M \implies \text{measure } M \ A > 0 \implies (1 / \text{measure } M \ A) *_R$ 
set-lebesgue-integral M A f  $\in S$ 
shows AE x in M. f x  $\in S$ 
proof (induct rule: sigma-finite-measure-induct)
case (finite-measure N  $\Omega$ )

```

interpret finite-measure N **by** (rule finite-measure)

```

have integrable[measurable]: integrable N f using assms finite-measure by (auto
simp: integrable-restrict-space integrable-mult-indicator)
have average:  $(1 / \text{Sigma-Algebra.measure } N \ A) *_R \text{set-lebesgue-integral } N \ A \ f$ 
 $\in S$  if  $A \in \text{sets } N$   $\text{measure } N \ A > 0$  for A
proof -
  have *:  $A \in \text{sets } M$  using that by (simp add: sets-restrict-space-iff finite-measure)
  have  $A = A \cap \Omega$  by (metis finite-measure(2,3) inf.orderE sets.sets-into-space
space-restrict-space that(1))
  hence set-lebesgue-integral N A f = set-lebesgue-integral M A f unfolding
finite-measure by (subst set-integral-restrict-space, auto simp add: finite-measure
set-lebesgue-integral-def indicator-inter-arith[symmetric])
  moreover have  $\text{measure } N \ A = \text{measure } M \ A$  using that by (auto intro!:
measure-restrict-space simp add: finite-measure sets-restrict-space-iff)
  ultimately show ?thesis using that * assms(3) by presburger
qed

```

obtain $D :: 'b$ set **where** $\text{balls-basis: topological-basis (case-prod ball ' (D \times (\mathbb{Q} \cap \{0 < ..\})))}$ **and** $\text{countable-D: countable D}$ **using** $\text{balls-countable-basis}$ **by** blast
have $\text{countable-balls: countable (case-prod ball ' (D \times (\mathbb{Q} \cap \{0 < ..\})))}$ **using** $\text{countable-rat countable-D}$ **by** blast

obtain B **where** $\text{B-balls: } B \subseteq \text{case-prod ball ' (D \times (\mathbb{Q} \cap \{0 < ..\}))} \cup B = -S$
using $\text{topological-basis[THEN iffD1, OF balls-basis] open-Compl[OF assms(2)]}$ **by** meson
hence $\text{countable-B: countable B}$ **using** $\text{countable-balls countable-subset}$ **by** fast

define b **where** $b = \text{from-nat-into (B \cup \{\{\}\})}$
have $B \cup \{\{\}\} \neq \{\}$ **by** simp
have $\text{range-b: range } b = B \cup \{\{\}\}$ **using** countable-B **by** $(\text{auto simp add: b-def intro!: range-from-nat-into})$
have $\text{open-b: open (b i) for } i$ **unfolding** $b\text{-def}$ **using** $\text{B-balls open-ball from-nat-into[of B \cup \{\{\}\} i]}$ **by** force
have $\text{Union-range-b: } \bigcup (\text{range } b) = -S$ **using** B-balls range-b **by** simp

{
fix $v \ r$ **assume** $\text{ball-in-Compl: ball } v \ r \subseteq -S$
define A **where** $A = f - ' \text{ball } v \ r \cap \text{space } N$
have $\text{dist-less: dist (f } x) \ v < r \text{ if } x \in A \text{ for } x$ **using** $\text{that unfolding A-def vimage-def}$ **by** $(\text{simp add: dist-commute})$
hence $\text{AE-less: AE } x \in A \text{ in } N. \text{ norm (f } x - v) < r$ **by** $(\text{auto simp add: dist-norm})$
have $*$: $A \in \text{sets } N$ **unfolding** $A\text{-def}$ **by** simp
have $\text{emeasure } N \ A = 0$
proof $-$
{
assume $\text{asm: emeasure } N \ A > 0$
hence $\text{measure-pos: measure } N \ A > 0$ **unfolding** $\text{emeasure-eq-measure}$ **by** simp
hence $(1 / \text{measure } N \ A) *_{\mathbb{R}} \text{set-lebesgue-integral } N \ A \ f - v = (1 / \text{measure } N \ A) *_{\mathbb{R}} \text{set-lebesgue-integral } N \ A \ (\lambda x. f \ x - v)$ **using** $\text{integrable integrable-const *}$ **by** $(\text{subst set-integral-diff(2), auto simp add: set-integrable-def set-integral-const[OF *] algebra-simps intro!: integrable-mult-indicator})$
moreover **have** $\text{norm } (\int x \in A. (f \ x - v) \partial N) \leq (\int x \in A. \text{norm (f } x - v) \partial N)$ **using** $*$ **by** $(\text{auto intro!: integral-norm-bound[of } N \ \lambda x. \text{indicator } A \ x *_{\mathbb{R}} (f \ x - v), \text{ THEN order-trans}] \text{integrable-mult-indicator integrable simp add: set-lebesgue-integral-def})$
ultimately **have** $\text{norm } ((1 / \text{measure } N \ A) *_{\mathbb{R}} \text{set-lebesgue-integral } N \ A \ f - v) \leq \text{set-lebesgue-integral } N \ A \ (\lambda x. \text{norm (f } x - v)) / \text{measure } N \ A$ **using** asm **by** $(\text{auto intro: divide-right-mono})$
also **have** $\dots < \text{set-lebesgue-integral } N \ A \ (\lambda x. r) / \text{measure } N \ A$
unfolding $\text{set-lebesgue-integral-def}$
using $\text{asm * integrable integrable-const AE-less measure-pos}$
by $(\text{intro divide-strict-right-mono integral-less-AE[of - - A] integrable-mult-indicator (fastforce simp add: dist-less dist-norm indicator-def)})+$

also have ... = r using * *measure-pos* by (*simp add: set-integral-const*)
 finally have $\text{dist } ((1 / \text{measure } N \ A) *_{\mathbb{R}} \text{set-lebesgue-integral } N \ A \ f) \ v < r$
 by (*subst dist-norm*)
 hence *False* using *average*[*OF * measure-pos*] by (*metis ComplD dist-commute in-mono mem-ball ball-in-Compl*)
 }
 thus ?thesis by *fastforce*
 qed
 }
 note * = *this*
 {
 fix $b' \in B$
 hence *ball-subset-Compl*: $b' \subseteq -S$ and *ball-radius-pos*: $\exists v \in D. \exists r > 0. b' = \text{ball } v \ r$ using *B-balls* by (*blast, fast*)
 }
 note ** = *this*
 hence $\text{emeasure } N \ (f -^{\cdot} b \ i \cap \text{space } N) = 0$ for i by (*cases b i = {}, simp*)
 (*metis UnE singletonD * range-b[THEN eq-refl, THEN range-subsetD]*)
 hence $\text{emeasure } N \ (\bigcup i. f -^{\cdot} b \ i \cap \text{space } N) = 0$ using *open-b* by (*intro emeasure-UN-eq-0 fastforce+*)
 moreover have $(\bigcup i. f -^{\cdot} b \ i \cap \text{space } N) = f -^{\cdot} (\bigcup (\text{range } b)) \cap \text{space } N$ by *blast*
 ultimately have $\text{emeasure } N \ (f -^{\cdot} (-S) \cap \text{space } N) = 0$ using *Union-range-b*
 by *argo*
 hence $AE \ x \text{ in } N. f \ x \notin -S$ using *open-Compl*[*OF assms(2)*] by (*intro AE-iff-measurable[THEN iffD2], auto*)
 thus ?case by *force*
 qed (*simp add: pred-sets2[OF borel-closed] assms(2)*)

lemma *density-nonneg*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
 assumes *integrable* $M \ f$
 and $\bigwedge A. A \in \text{sets } M \implies \text{set-lebesgue-integral } M \ A \ f \geq 0$
 shows $AE \ x \text{ in } M. f \ x \geq 0$
 using *averaging-theorem*[*OF assms(1), of {0..}, OF closed-atLeast*] *assms(2)*
 by (*simp add: scaleR-nonneg-nonneg*)

lemma *density-zero*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
 assumes *integrable* $M \ f$
 and *density-0*: $\bigwedge A. A \in \text{sets } M \implies \text{set-lebesgue-integral } M \ A \ f = 0$
 shows $AE \ x \text{ in } M. f \ x = 0$
 using *averaging-theorem*[*OF assms(1), of {0}*] *assms(2)*
 by (*simp add: scaleR-nonneg-nonneg*)

lemma *density-unique*:

fixes $f \ f' :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
 assumes *integrable* $M \ f$ *integrable* $M \ f'$
 and *density-eq*: $\bigwedge A. A \in \text{sets } M \implies \text{set-lebesgue-integral } M \ A \ f = \text{set-lebesgue-integral}$


```

M A f'
shows AE x in M. f x = f' x
proof -
{
  fix A assume asm: A ∈ sets M
  hence LINT x|M. indicat-real A x *R (f x - f' x) = 0 using density-eq
  assms(1,2) by (simp add: set-lebesgue-integral-def algebra-simps Bochner-Integration.integral-diff[OF
  integrable-mult-indicator(1,1)])
}
thus ?thesis using density-zero[OF Bochner-Integration.integrable-diff[OF assms(1,2)]]
by (simp add: set-lebesgue-integral-def)
qed

```

```

lemma integral-nonneg-AE-eq-0-iff-AE:
  fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
  dered-real-vector}
  assumes f[measurable]: integrable M f and nonneg: AE x in M. 0 ≤ f x
  shows integralL M f = 0 ⟷ (AE x in M. f x = 0)
proof
  assume *: integralL M f = 0
  {
    fix A assume asm: A ∈ sets M
    have 0 ≤ integralL M (λx. indicator A x *R f x) using nonneg by (subst inte-
    gral-zero[of M, symmetric], intro integral-mono-AE-banach integrable-mult-indicator
    asm f integrable-zero, auto simp add: indicator-def)
    moreover have ... ≤ integralL M f using nonneg by (intro integral-mono-AE-banach
    integrable-mult-indicator asm f, auto simp add: indicator-def)
    ultimately have set-lebesgue-integral M A f = 0 unfolding set-lebesgue-integral-def
    using * by force
  }
  thus AE x in M. f x = 0 by (intro density-zero f, blast)
qed (auto simp add: integral-eq-zero-AE)

```

```

lemma integral-eq-mono-AE-eq-AE:
  fixes f g :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
  dered-real-vector}
  assumes integrable M f integrable M g integralL M f = integralL M g AE x in
  M. f x ≤ g x
  shows AE x in M. f x = g x
proof -
  define h where h = (λx. g x - f x)
  have AE x in M. h x = 0 unfolding h-def using assms by (subst inte-
  gral-nonneg-AE-eq-0-iff-AE[symmetric]) auto
  then show ?thesis unfolding h-def by auto
qed

```

end

end

```

theory Filtration
imports HOL-Probability.Conditional-Expectation HOL-Probability.Stopping-Time
Measure-Space-Addendum
begin

```

1.3 Filtered Sigma Finite Measure

```

locale filtered-sigma-finite-measure = sigma-finite-measure M + filtration space M
for M and F :: 't :: {second-countable-topology, linorder-topology, order-bot}  $\Rightarrow$ 
'a measure +
  assumes subalgebra:  $\bigwedge i. \text{subalgebra } M (F i)$ 
  and sigma-finite: sigma-finite-measure (restr-to-subalg M (F bot))

```

```

locale ennreal-filtered-sigma-finite-measure = filtered-sigma-finite-measure M F for
M and F :: ennreal  $\Rightarrow$  -
locale nat-filtered-sigma-finite-measure = filtered-sigma-finite-measure M F for M
and F :: nat  $\Rightarrow$  -

```

```

sublocale filtered-sigma-finite-measure  $\subseteq$  sigma-finite-subalgebra M F i by (metis
bot.extremum sigma-finite sigma-finite-subalgebra.intro subalgebra sets-F-mono sigma-finite-subalgebra.nested-s
subalgebra-def)

```

1.4 Natural Filtration

```

definition natural-filtration :: 'a measure  $\Rightarrow$  's measure  $\Rightarrow$  ('t  $\Rightarrow$  'a  $\Rightarrow$  's)  $\Rightarrow$  't ::
{second-countable-topology, linorder-topology, order-bot}  $\Rightarrow$  'a measure where
  natural-filtration M N Y = ( $\lambda t. \text{restr-to-subalg } M (\text{sigma-gen } (\text{space } M) N \{Y i \mid i. i \leq t\})$ )

```

```

lemma
  assumes  $\bigwedge i. Y i \in M \rightarrow_M N$ 
  shows sets-natural-filtration[simp]: sets (natural-filtration M N Y t) = sigma-sets
(space M) ( $\bigcup i \leq t. \{Y i - 'A \cap \text{space } M \mid A. A \in N\}$ )
  and space-natural-filtration[simp]: space (natural-filtration M N Y t) = space M
  by (standard; (subst natural-filtration-def, subst sets-restr-to-subalg)) (auto simp
add: natural-filtration-def space-restr-to-subalg subalgebra-def intro!: sets.sigma-sets-subset
measurable-sets[OF assms] sigma-sets-mono)

```

```

end
theory Conditional-Expectation-Banach
imports HOL-Probability.Conditional-Expectation Sigma-Finite-Measure-Addendum
Bochner-Integration-Addendum Elementary-Metric-Spaces-Addendum
begin

```

```

definition has-cond-exp :: 'a measure  $\Rightarrow$  'a measure  $\Rightarrow$  ('a  $\Rightarrow$  'b)  $\Rightarrow$  ('a  $\Rightarrow$  'b::{real-normed-vector,
second-countable-topology})  $\Rightarrow$  bool where
  has-cond-exp M F f g = ( $(\forall A \in \text{sets } F. (\int x \in A. f x \, \partial M) = (\int x \in A. g x \, \partial M))$ 
 $\wedge$  integrable M f
 $\wedge$  integrable M g

```

$\wedge g \in \text{borel-measurable } F)$

lemma *has-cond-expI'*[intro]:

assumes $\bigwedge A. A \in \text{sets } F \implies (\int x \in A. f \ x \ \partial M) = (\int x \in A. g \ x \ \partial M)$
 $\text{integrable } M \ f$
 $\text{integrable } M \ g$
 $g \in \text{borel-measurable } F$
shows *has-cond-exp* $M \ F \ f \ g$
using *assms* **unfolding** *has-cond-exp-def* **by** *simp*

lemma *has-cond-expD*:

assumes *has-cond-exp* $M \ F \ f \ g$
shows $\bigwedge A. A \in \text{sets } F \implies (\int x \in A. f \ x \ \partial M) = (\int x \in A. g \ x \ \partial M)$
 $\text{integrable } M \ f$
 $\text{integrable } M \ g$
 $g \in \text{borel-measurable } F$
using *assms* **unfolding** *has-cond-exp-def* **by** *simp+*

definition *cond-exp* :: $'a \text{ measure} \Rightarrow 'a \text{ measure} \Rightarrow ('a \Rightarrow 'b) \Rightarrow ('a \Rightarrow 'b::\{\text{banach, second-countable-topology}\})$ **where**

cond-exp $M \ F \ f = (\text{if } \exists g. \text{has-cond-exp } M \ F \ f \ g \text{ then } (\text{SOME } g. \text{has-cond-exp } M \ F \ f \ g) \text{ else } (\lambda \cdot. 0))$

lemma *borel-measurable-cond-exp*[*measurable*]: *cond-exp* $M \ F \ f \in \text{borel-measurable } F$

by (*metis cond-exp-def someI has-cond-exp-def borel-measurable-const*)

lemma *integrable-cond-exp*[intro]: *integrable* $M \ (\text{cond-exp } M \ F \ f)$

by (*metis cond-exp-def has-cond-expD(3) integrable-zero someI*)

lemma *set-integrable-cond-exp*[intro]:

assumes $A \in \text{sets } M$

shows *set-integrable* $M \ A \ (\text{cond-exp } M \ F \ f)$ **using** *integrable-mult-indicator*[*OF assms integrable-cond-exp, of F f*] **by** (*auto simp add: set-integrable-def intro!: integrable-mult-indicator*[*OF assms integrable-cond-exp*])

context *sigma-finite-subalgebra*

begin

lemma *borel-measurable-cond-exp'*[*measurable*]: *cond-exp* $M \ F \ f \in \text{borel-measurable } M$

by (*metis cond-exp-def someI has-cond-exp-def borel-measurable-const subalg measurable-from-subalg*)

lemma *cond-exp-null*:

assumes $\nexists g. \text{has-cond-exp } M \ F \ f \ g$

shows *cond-exp* $M \ F \ f = (\lambda \cdot. 0)$

unfolding *cond-exp-def* **using** *assms* **by** *argo*

lemma *has-cond-exp-charact*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes *has-cond-exp* $M F f g$
shows *has-cond-exp* $M F f$ (*cond-exp* $M F f$)
 $AE x \text{ in } M. \text{cond-exp } M F f x = g x$
proof –
show *cond-exp*: *has-cond-exp* $M F f$ (*cond-exp* $M F f$) **using** *assms someI*
cond-exp-def **by** *metis*
let $?MF = \text{restr-to-subalg } M F$
interpret *sigma-finite-measure* $?MF$ **by** (*rule sigma-fin-subalg*)
{
fix A **assume** $A \in \text{sets } ?MF$
then have [*measurable*]: $A \in \text{sets } F$ **using** *sets-restr-to-subalg[OF subalg]* **by**
simp
have $(\int x \in A. g x \partial ?MF) = (\int x \in A. g x \partial M)$ **using** *assms subalg* **by** (*auto*
simp add: integral-subalgebra2 set-lebesgue-integral-def dest!: has-cond-expD)
also have $\dots = (\int x \in A. \text{cond-exp } M F f x \partial M)$ **using** *assms cond-exp* **by**
(*simp add: has-cond-exp-def*)
also have $\dots = (\int x \in A. \text{cond-exp } M F f x \partial ?MF)$ **using** *subalg* **by** (*auto simp*
add: integral-subalgebra2 set-lebesgue-integral-def)
finally have $(\int x \in A. g x \partial ?MF) = (\int x \in A. \text{cond-exp } M F f x \partial ?MF)$ **by**
simp
}
hence $AE x \text{ in } ?MF. \text{cond-exp } M F f x = g x$ **using** *cond-exp assms subalg* **by**
(*intro density-unique, auto dest: has-cond-expD intro!: integrable-in-subalg*)
then show $AE x \text{ in } M. \text{cond-exp } M F f x = g x$ **using** *AE-restr-to-subalg[OF*
subalg] **by** *simp*
qed

lemma *cond-exp-F-meas[intro, simp]*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes *integrable* $M f$
 $f \in \text{borel-measurable } F$
shows $AE x \text{ in } M. \text{cond-exp } M F f x = f x$
by (*rule has-cond-exp-charact(2), auto intro: assms*)

Congruence

lemma *has-cond-exp-cong*:
assumes *integrable* $M f \wedge x. x \in \text{space } M \implies f x = g x$ *has-cond-exp* $M F g h$
shows *has-cond-exp* $M F f h$
proof (*intro has-cond-expI'[OF - assms(1)], goal-cases*)
case (1 A)
hence *set-lebesgue-integral* $M A f = \text{set-lebesgue-integral } M A g$ **by** (*intro set-lebesgue-integral-cong*)
(*meson assms(2) subalg in-mono subalgebra-def sets.sets-into-space subalgebra-def*
subsetD) +
then show $?case$ **using** 1 *assms(3)* **by** (*simp add: has-cond-exp-def*)
qed (*auto simp add: has-cond-expD[OF assms(3)]*)

lemma *cond-exp-cong*:

```

fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
assumes  $\text{integrable } M f \text{ integrable } M g \wedge x. x \in \text{space } M \implies f x = g x$ 
shows  $AE\ x \text{ in } M. \text{cond-exp } M F f x = \text{cond-exp } M F g x$ 
proof (cases  $\exists h. \text{has-cond-exp } M F f h$ )
  case True
    then obtain  $h$  where  $h: \text{has-cond-exp } M F f h \text{ has-cond-exp } M F g h$  using
 $\text{has-cond-exp-cong assms}$  by metis
    show ?thesis using  $h[\text{THEN } \text{has-cond-exp-charact}(2)]$  by fastforce
  next
    case False
    moreover have  $\nexists h. \text{has-cond-exp } M F g h$  using False has-cond-exp-cong assms
by auto
    ultimately show ?thesis unfolding cond-exp-def by auto
qed

```

```

lemma has-cond-exp-cong-AE:
assumes  $\text{integrable } M f \text{ AE } x \text{ in } M. f x = g x \text{ has-cond-exp } M F g h$ 
shows  $\text{has-cond-exp } M F f h$ 
using  $\text{assms}(1, 2) \text{ subalg subalgebra-def subset-iff}$ 
by (intro  $\text{has-cond-expI'}$ , subst set-lebesgue-integral-cong-AE[OF -  $\text{assms}(1)[\text{THEN } \text{borel-measurable-integrable}] \text{ borel-measurable-integrable}(1)[\text{OF } \text{has-cond-expD}(2)[\text{OF } \text{assms}(3)]]]$ ])
  (fast intro:  $\text{has-cond-expD}[\text{OF } \text{assms}(3)] \text{ integrable-cong-AE-imp}[\text{OF} - - \text{AE-symmetric}]) +$ 

```

```

lemma has-cond-exp-cong-AE':
assumes  $h \in \text{borel-measurable } F \text{ AE } x \text{ in } M. h x = h' x \text{ has-cond-exp } M F f h'$ 
shows  $\text{has-cond-exp } M F f h$ 
using  $\text{assms}(1, 2) \text{ subalg subalgebra-def subset-iff}$ 
using  $\text{AE-restr-to-subalg2}[\text{OF } \text{subalg } \text{assms}(2)] \text{ measurable-from-subalg}$ 
by (intro  $\text{has-cond-expI'}$ , subst set-lebesgue-integral-cong-AE[OF -  $\text{measurable-from-subalg}(1, 1)[\text{OF } \text{subalg}], \text{OF} - \text{assms}(1) \text{ has-cond-expD}(4)[\text{OF } \text{assms}(3)]]]$ ])
  (fast intro:  $\text{has-cond-expD}[\text{OF } \text{assms}(3)] \text{ integrable-cong-AE-imp}[\text{OF} - - \text{AE-symmetric}]) +$ 

```

```

lemma cond-exp-cong-AE:
fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
assumes  $\text{integrable } M f \text{ integrable } M g \text{ AE } x \text{ in } M. f x = g x$ 
shows  $AE\ x \text{ in } M. \text{cond-exp } M F f x = \text{cond-exp } M F g x$ 
proof (cases  $\exists h. \text{has-cond-exp } M F f h$ )
  case True
    then obtain  $h$  where  $h: \text{has-cond-exp } M F f h \text{ has-cond-exp } M F g h$  using
 $\text{has-cond-exp-cong-AE assms}$  by (metis (mono-tags, lifting) eventually-mono)
    show ?thesis using  $h[\text{THEN } \text{has-cond-exp-charact}(2)]$  by fastforce
  next
    case False
    moreover have  $\nexists h. \text{has-cond-exp } M F g h$  using False has-cond-exp-cong-AE assms by auto
    ultimately show ?thesis unfolding cond-exp-def by auto
qed

```

lemma *has-cond-exp-real*:

fixes $f :: 'a \Rightarrow \text{real}$

assumes *integrable* $M f$

shows *has-cond-exp* $M F f$ (*real-cond-exp* $M F f$)

by (*standard*, *auto intro!*: *real-cond-exp-intA assms*)

lemma *cond-exp-real[intro]*:

fixes $f :: 'a \Rightarrow \text{real}$

assumes *integrable* $M f$

shows $AE\ x\ in\ M. \text{cond-exp}\ M F f\ x = \text{real-cond-exp}\ M F f\ x$

using *has-cond-exp-charact has-cond-exp-real assms* **by** *blast*

lemma *cond-exp-cmult*:

fixes $f :: 'a \Rightarrow \text{real}$

assumes *integrable* $M f$

shows $AE\ x\ in\ M. \text{cond-exp}\ M F (\lambda x. c * f\ x)\ x = c * \text{cond-exp}\ M F f\ x$

using *real-cond-exp-cmult[OF assms(1), of c] assms(1)[THEN cond-exp-real] assms(1)[THEN integrable-mult-right, THEN cond-exp-real, of c]* **by** *fastforce*

Indicator functions

lemma *has-cond-exp-indicator*:

assumes $A \in \text{sets}\ M\ \text{emeasure}\ M\ A < \infty$

shows *has-cond-exp* $M F (\lambda x. \text{indicat-real}\ A\ x *_{\mathbb{R}} y)$ ($\lambda x. \text{real-cond-exp}\ M F (\text{indicator}\ A)\ x *_{\mathbb{R}} y$)

proof (*intro has-cond-expI', goal-cases*)

case ($1\ B$)

have $\int_{x \in B. (\text{indicat-real}\ A\ x *_{\mathbb{R}} y)\ \partial M} = (\int_{x \in B. \text{indicat-real}\ A\ x\ \partial M}) *_{\mathbb{R}} y$ **using** *assms* **by** (*intro set-integral-scaleR-left, meson 1 in-mono subalg subalgebra-def, blast*)

also have $\dots = (\int_{x \in B. \text{real-cond-exp}\ M F (\text{indicator}\ A)\ x\ \partial M}) *_{\mathbb{R}} y$ **using** 1 *assms* **by** (*subst real-cond-exp-intA, auto*)

also have $\dots = \int_{x \in B. (\text{real-cond-exp}\ M F (\text{indicator}\ A)\ x *_{\mathbb{R}} y)\ \partial M}$ **using** *assms* **by** (*intro set-integral-scaleR-left[symmetric], meson 1 in-mono subalg subalgebra-def, blast*)

finally show *?case* .

next

case 2

then show *?case* **using** *integrable-scaleR-left integrable-real-indicator assms* **by** *blast*

next

case 3

show *?case* **using** *assms* **by** (*intro integrable-scaleR-left, intro real-cond-exp-int, blast+*)

next

case 4

then show *?case* **by** (*intro borel-measurable-scaleR, intro Conditional-Expectation.borel-measurable-cond-exp, simp*)

qed

lemma *cond-exp-indicator*[intro]:
fixes $y :: 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes [measurable]: $A \in \text{sets } M \text{ emeasure } M A < \infty$
shows $\text{AE } x \text{ in } M. \text{ cond-exp } M F (\lambda x. \text{ indicat-real } A x *_R y) x = \text{ cond-exp } M F (\text{indicator } A) x *_R y$
proof –
have $\text{AE } x \text{ in } M. \text{ cond-exp } M F (\lambda x. \text{ indicat-real } A x *_R y) x = \text{ real-cond-exp } M F (\text{indicator } A) x *_R y$ **using** *has-cond-exp-indicator*[OF *assms*] *has-cond-exp-charact*
by *blast*
thus ?thesis **using** *cond-exp-real*[OF *integrable-real-indicator*, OF *assms*] **by** *fast-force*
qed

Addition

lemma *has-cond-exp-add*:
fixes $f g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *has-cond-exp* $M F f f'$ *has-cond-exp* $M F g g'$
shows *has-cond-exp* $M F (\lambda x. f x + g x) (\lambda x. f' x + g' x)$
proof (intro *has-cond-expI'*, goal-cases)
case (1 A)
have $\int_{x \in A}. (f x + g x) \partial M = (\int_{x \in A}. f x \partial M) + (\int_{x \in A}. g x \partial M)$ **using** *assms*[*THEN* *has-cond-expD*(2)] *subalg* 1 **by** (intro *set-integral-add*(2), *auto simp add: subalgebra-def set-integrable-def* intro: *integrable-mult-indicator*)
also **have** $\dots = (\int_{x \in A}. f' x \partial M) + (\int_{x \in A}. g' x \partial M)$ **using** *assms*[*THEN* *has-cond-expD*(1)[OF - 1]] **by** *arg0*
also **have** $\dots = \int_{x \in A}. (f' x + g' x) \partial M$ **using** *assms*[*THEN* *has-cond-expD*(3)] *subalg* 1 **by** (intro *set-integral-add*(2)[*symmetric*], *auto simp add: subalgebra-def set-integrable-def* intro: *integrable-mult-indicator*)
finally **show** ?case .
next
case 2
then **show** ?case **by** (metis *Bochner-Integration.integrable-add* *assms* *has-cond-expD*(2))
next
case 3
then **show** ?case **by** (metis *Bochner-Integration.integrable-add* *assms* *has-cond-expD*(3))
next
case 4
then **show** ?case **using** *assms* *borel-measurable-add* *has-cond-expD*(4) **by** *blast*
qed

lemma *has-cond-exp-scaleR-right*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *has-cond-exp* $M F f f'$
shows *has-cond-exp* $M F (\lambda x. c *_R f x) (\lambda x. c *_R f' x)$
using *has-cond-expD*[OF *assms*] **by** (intro *has-cond-expI'*, *auto*)

lemma *cond-exp-scaleR-right*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f$

```

  shows  $\text{AE } x \text{ in } M. \text{ cond-exp } M F (\lambda x. c *_R f x) x = c *_R \text{ cond-exp } M F f x$ 
proof (cases  $\exists f'. \text{ has-cond-exp } M F f f'$ )
  case True
  then show ?thesis using assms has-cond-exp-charact has-cond-exp-scaleR-right
  by metis
next
  case False
  show ?thesis
  proof (cases  $c = 0$ )
    case True
    then show ?thesis by simp
  next
    case c-nonzero: False
    have  $\nexists f'. \text{ has-cond-exp } M F (\lambda x. c *_R f x) f'$ 
    proof (standard, goal-cases)
      case 1
      then obtain  $f'$  where  $f': \text{ has-cond-exp } M F (\lambda x. c *_R f x) f'$  by blast
      have  $\text{ has-cond-exp } M F f (\lambda x. \text{ inverse } c *_R f' x)$  using has-cond-expD[OF  $f'$ ]
        divideR-right[OF c-nonzero] assms by (intro has-cond-expI', auto)
      then show ?case using False by blast
    qed
    then show ?thesis using cond-exp-null[OF False] cond-exp-null by force
  qed
qed

```

```

lemma cond-exp-uminus:
  fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
  assumes integrable  $M f$ 
  shows  $\text{AE } x \text{ in } M. \text{ cond-exp } M F (\lambda x. - f x) x = - \text{ cond-exp } M F f x$ 
  using cond-exp-scaleR-right[OF assms, of  $-1$ ] by force

```

```

lemma has-cond-exp-simple:
  fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
  assumes simple-function  $M f$  emeasure  $M \{y \in \text{space } M. f y \neq 0\} \neq \infty$ 
  shows  $\text{ has-cond-exp } M F f (\text{ cond-exp } M F f)$ 
  using assms
proof (induction rule: simple-integrable-function-induct)
  case (cong  $f g$ )
  then show ?case using has-cond-exp-cong by (metis (no-types, opaque-lifting)
    Bochner-Integration.integrable-cong has-cond-expD(2) has-cond-exp-charact(1))
next
  case (indicator  $A y$ )
  then show ?case using has-cond-exp-charact[OF has-cond-exp-indicator] by fast
next
  case (add  $u v$ )
  then show ?case using has-cond-exp-add has-cond-exp-charact(1) by blast
qed

```

```

lemma cond-exp-contraction-real:

```


fixes $f :: 'a \Rightarrow \text{real}$
assumes $\text{integrable}[\text{measurable}]$: $\text{integrable } M f$
shows $\text{AE } x \text{ in } M. \text{norm } (\text{cond-exp } M F f x) \leq \text{cond-exp } M F (\lambda x. \text{norm } (f x)) x$
proof –
have int : $\text{integrable } M (\lambda x. \text{norm } (f x))$ **using** assms **by** blast
have *: $\text{AE } x \text{ in } M. 0 \leq \text{cond-exp } M F (\lambda x. \text{norm } (f x)) x$ **using** $\text{cond-exp-real}[\text{THEN } \text{AE-symmetric}, \text{OF integrable-norm}[\text{OF integrable}]] \text{real-cond-exp-ge-c}[\text{OF integrable-norm}[\text{OF integrable}], \text{of } 0] \text{norm-ge-zero}$ **by** fastforce
have **: $A \in \text{sets } F \implies \int x \in A. |f x| \partial M = \int x \in A. \text{real-cond-exp } M F (\lambda x. \text{norm } (f x)) x \partial M$ **for** A **unfolding** real-norm-def **using** assms integrable-abs $\text{real-cond-exp-intA}$ **by** blast

have norm-int : $A \in \text{sets } F \implies (\int x \in A. |f x| \partial M) = (\int^+ x \in A. |f x| \partial M)$ **for** A **using** assms **by** $(\text{intro nn-set-integral-eq-set-integral}[\text{symmetric}], \text{blast}, \text{fastforce})$ $(\text{meson subalg subalgebra-def subsetD})$

have $\text{AE } x \text{ in } M. \text{real-cond-exp } M F (\lambda x. \text{norm } (f x)) x \geq 0$ **using** $\text{int real-cond-exp-ge-c}$ **by** force
hence cond-exp-norm-int : $A \in \text{sets } F \implies (\int x \in A. \text{real-cond-exp } M F (\lambda x. \text{norm } (f x)) x \partial M) = (\int^+ x \in A. \text{real-cond-exp } M F (\lambda x. \text{norm } (f x)) x \partial M)$ **for** A **using** assms **by** $(\text{intro nn-set-integral-eq-set-integral}[\text{symmetric}], \text{blast}, \text{fastforce})$ $(\text{meson subalg subalgebra-def subsetD})$

have $A \in \text{sets } F \implies \int^+ x \in A. |f x| \partial M = \int^+ x \in A. \text{real-cond-exp } M F (\lambda x. \text{norm } (f x)) x \partial M$ **for** A **using** ** $\text{norm-int cond-exp-norm-int}$ **by** $(\text{auto simp add: nn-integral-set-ennreal})$
moreover **have** $(\lambda x. \text{ennreal } |f x|) \in \text{borel-measurable } M$ **by** measurable
moreover **have** $(\lambda x. \text{ennreal } (\text{real-cond-exp } M F (\lambda x. \text{norm } (f x)) x)) \in \text{borel-measurable } F$ **by** measurable
ultimately **have** $\text{AE } x \text{ in } M. \text{nn-cond-exp } M F (\lambda x. \text{ennreal } |f x|) x = \text{real-cond-exp } M F (\lambda x. \text{norm } (f x)) x$ **by** $(\text{intro nn-cond-exp-charact}[\text{THEN AE-symmetric}], \text{auto})$
hence $\text{AE } x \text{ in } M. \text{nn-cond-exp } M F (\lambda x. \text{ennreal } |f x|) x \leq \text{cond-exp } M F (\lambda x. \text{norm } (f x)) x$ **using** $\text{cond-exp-real}[\text{OF int}]$ **by** force
moreover **have** $\text{AE } x \text{ in } M. |\text{real-cond-exp } M F f x| = \text{norm } (\text{cond-exp } M F f x)$
unfolding real-norm-def **using** $\text{cond-exp-real}[\text{OF assms}]$ * **by** force
ultimately **have** $\text{AE } x \text{ in } M. \text{ennreal } (\text{norm } (\text{cond-exp } M F f x)) \leq \text{cond-exp } M F (\lambda x. \text{norm } (f x)) x$ **using** $\text{real-cond-exp-abs}[\text{OF assms}[\text{THEN borel-measurable-integrable}]]$ **by** fastforce
hence $\text{AE } x \text{ in } M. \text{enn2real } (\text{ennreal } (\text{norm } (\text{cond-exp } M F f x))) \leq \text{enn2real } (\text{cond-exp } M F (\lambda x. \text{norm } (f x)) x)$ **using** ennreal-le-iff2 **by** force
thus $?thesis$ **using** * **by** fastforce
qed

lemma $\text{cond-exp-contraction-simple}$:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes $\text{simple-function } M f \text{emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty$
shows $\text{AE } x \text{ in } M. \text{norm } (\text{cond-exp } M F f x) \leq \text{cond-exp } M F (\lambda x. \text{norm } (f x)) x$
using assms

proof (*induction rule: simple-integrable-function-induct*)
case (*cong f g*)
hence $ae: AE\ x\ in\ M. f\ x = g\ x$ **by** *blast*
hence $AE\ x\ in\ M. cond-exp\ M\ F\ f\ x = cond-exp\ M\ F\ g\ x$ **using** *cong has-cond-exp-simple*
by (*subst cond-exp-cong-AE*) (*auto intro!: has-cond-expD(2)*)
hence $AE\ x\ in\ M. norm\ (cond-exp\ M\ F\ f\ x) = norm\ (cond-exp\ M\ F\ g\ x)$ **by**
force
moreover **have** $AE\ x\ in\ M. cond-exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x = cond-exp\ M\ F\ (\lambda x. norm\ (g\ x))\ x$ **using** *ae cong has-cond-exp-simple* **by** (*subst cond-exp-cong-AE*)
(*auto dest: has-cond-expD*)
ultimately show *?case using cong(6)* **by** *fastforce*
next
case (*indicator A y*)
hence $AE\ x\ in\ M. cond-exp\ M\ F\ (\lambda a. indicator\ A\ a\ *_R\ y)\ x = cond-exp\ M\ F\ (indicator\ A)\ x\ *_R\ y$ **by** *blast*
hence $*$: $AE\ x\ in\ M. norm\ (cond-exp\ M\ F\ (\lambda a. indicat-real\ A\ a\ *_R\ y)\ x) \leq norm\ y$
 $\ast cond-exp\ M\ F\ (\lambda x. norm\ (indicat-real\ A\ x))\ x$ **using** *cond-exp-contraction-real[OF integrable-real-indicator, OF indicator]* **by** *fastforce*

have $AE\ x\ in\ M. norm\ y\ \ast\ cond-exp\ M\ F\ (\lambda x. norm\ (indicat-real\ A\ x))\ x = norm\ y$
 $\ast\ real-cond-exp\ M\ F\ (\lambda x. norm\ (indicat-real\ A\ x))\ x$ **using** *cond-exp-real[OF integrable-real-indicator, OF indicator]* **by** *fastforce*
moreover **have** $AE\ x\ in\ M. cond-exp\ M\ F\ (\lambda x. norm\ y\ \ast\ norm\ (indicat-real\ A\ x))\ x =$
 $real-cond-exp\ M\ F\ (\lambda x. norm\ y\ \ast\ norm\ (indicat-real\ A\ x))\ x$ **using**
indicator **by** (*intro cond-exp-real, auto*)
ultimately have $AE\ x\ in\ M. norm\ y\ \ast\ cond-exp\ M\ F\ (\lambda x. norm\ (indicat-real\ A\ x))\ x =$
 $cond-exp\ M\ F\ (\lambda x. norm\ y\ \ast\ norm\ (indicat-real\ A\ x))\ x$ **using** *real-cond-exp-cmult[of*
 $\lambda x. norm\ (indicat-real\ A\ x)\ norm\ y]$ *indicator* **by** *fastforce*
moreover **have** $(\lambda x. norm\ y\ \ast\ norm\ (indicat-real\ A\ x)) = (\lambda x. norm\ (indicat-real\ A\ x\ *_R\ y))$
by *force*
ultimately show *?case using ** **by** *force*
next
case (*add u v*)
have $AE\ x\ in\ M. norm\ (cond-exp\ M\ F\ (\lambda a. u\ a + v\ a)\ x) = norm\ (cond-exp\ M\ F\ u\ x +$
 $cond-exp\ M\ F\ v\ x)$ **using** *has-cond-exp-charact(2)[OF has-cond-exp-add, OF has-cond-exp-simple(1,1), OF add(1,2,3,4)]* **by** *fastforce*
moreover **have** $AE\ x\ in\ M. norm\ (cond-exp\ M\ F\ u\ x + cond-exp\ M\ F\ v\ x) \leq$
 $norm\ (cond-exp\ M\ F\ u\ x) + norm\ (cond-exp\ M\ F\ v\ x)$ **using** *norm-triangle-ineq*
by *blast*
moreover **have** $AE\ x\ in\ M. norm\ (cond-exp\ M\ F\ u\ x) + norm\ (cond-exp\ M\ F\ v\ x) \leq$
 $cond-exp\ M\ F\ (\lambda x. norm\ (u\ x))\ x + cond-exp\ M\ F\ (\lambda x. norm\ (v\ x))\ x$ **using**
add(6,7) **by** *fastforce*
moreover **have** $AE\ x\ in\ M. cond-exp\ M\ F\ (\lambda x. norm\ (u\ x))\ x + cond-exp\ M\ F\ (\lambda x. norm\ (v\ x))\ x =$
 $cond-exp\ M\ F\ (\lambda x. norm\ (u\ x) + norm\ (v\ x))\ x$ **using** *integrable-simple-function[OF add(1,2)]*
integrable-simple-function[OF add(3,4)] **by**
(*intro has-cond-exp-charact(2)[OF has-cond-exp-add[OF has-cond-exp-charact(1,1)], THEN AE-symmetric], auto intro: has-cond-exp-real*)
moreover **have** $AE\ x\ in\ M. cond-exp\ M\ F\ (\lambda x. norm\ (u\ x) + norm\ (v\ x))\ x =$
 $cond-exp\ M\ F\ (\lambda x. norm\ (u\ x + v\ x))\ x$ **using** *add(5) integrable-simple-function[OF*

add(1,2)] *integrable-simple-function*[*OF add*(3,4)] **by** (*intro cond-exp-cong*, *auto*)
ultimately show ?*case* **by** *force*
qed

lemma *has-cond-exp-lim*:

fixes *f* :: 'a \Rightarrow 'b::{*second-countable-topology*, *banach*}
assumes *integrable*[*measurable*]: *integrable* *M* *f*
and $\bigwedge i$. *simple-function* *M* (*s* *i*)
and $\bigwedge i$. *emeasure* *M* {*y* \in *space* *M*. *s* *i* *y* \neq 0} $\neq \infty$
and $\bigwedge x$. *x* \in *space* *M* \implies (λi . *s* *i* *x*) \longrightarrow *f* *x*
and $\bigwedge x$ *i*. *x* \in *space* *M* \implies *norm* (*s* *i* *x*) $\leq 2 * \text{norm } (f\ x)$
obtains *r*
where *has-cond-exp* *M* *F* *f* (λx . *lim* (λi . *cond-exp* *M* *F* (*s* (*r* *i*)) *x*))
AE *x* *in* *M*. *convergent* (λi . *cond-exp* *M* *F* (*s* (*r* *i*)) *x*)
strict-mono *r*

proof –

have [*measurable*]: (*s* *i*) \in *borel-measurable* *M* **for** *i* **using** *assms*(2) **by** (*simp* *add*: *borel-measurable-simple-function*)
have *integrable-s*: *integrable* *M* (λx . *s* *i* *x*) **for** *i* **using** *assms*(2) *assms*(3) *integrable-simple-function* **by** *blast*
have *integrable-4f*: *integrable* *M* (λx . 4 * *norm* (*f* *x*)) **using** *assms*(1) **by** *simp*
have *integrable-2f*: *integrable* *M* (λx . 2 * *norm* (*f* *x*)) **using** *assms*(1) **by** *simp*
have *integrable-2-cond-exp-norm-f*: *integrable* *M* (λx . 2 * *cond-exp* *M* *F* (λx . *norm* (*f* *x*)) *x*) **by** *fast*

have *emeasure* *M* {*y* \in *space* *M*. *s* *i* *y* – *s* *j* *y* \neq 0} \leq *emeasure* *M* {*y* \in *space* *M*. *s* *i* *y* \neq 0} + *emeasure* *M* {*y* \in *space* *M*. *s* *j* *y* \neq 0} **for** *i* *j* **using** *simple-functionD*(2)[*OF assms*(2)] **by** (*intro order-trans*[*OF emeasure-mono emeasure-subadditive*], *auto*)

hence *fin-sup*: *emeasure* *M* {*y* \in *space* *M*. *s* *i* *y* – *s* *j* *y* \neq 0} $\neq \infty$ **for** *i* *j* **using** *assms*(3) **by** (*metis* (*mono-tags*) *ennreal-add-eq-top* *linorder-not-less-top.not-eq-extremum* *infinity-ennreal-def*)

have *emeasure* *M* {*y* \in *space* *M*. *norm* (*s* *i* *y* – *s* *j* *y*) \neq 0} \leq *emeasure* *M* {*y* \in *space* *M*. *s* *i* *y* \neq 0} + *emeasure* *M* {*y* \in *space* *M*. *s* *j* *y* \neq 0} **for** *i* *j* **using** *simple-functionD*(2)[*OF assms*(2)] **by** (*intro order-trans*[*OF emeasure-mono emeasure-subadditive*], *auto*)

hence *fin-sup-norm*: *emeasure* *M* {*y* \in *space* *M*. *norm* (*s* *i* *y* – *s* *j* *y*) \neq 0} $\neq \infty$ **for** *i* *j* **using** *assms*(3) **by** (*metis* (*mono-tags*) *ennreal-add-eq-top* *linorder-not-less-top.not-eq-extremum* *infinity-ennreal-def*)

have *Cauchy*: *Cauchy* (λn . *s* *n* *x*) **if** *x* \in *space* *M* **for** *x* **using** *assms*(4) *LIM-SEQ-imp-Cauchy* **that** **by** *blast*

hence *bounded-range-s*: *bounded* (*range* (λn . *s* *n* *x*)) **if** *x* \in *space* *M* **for** *x* **using** *that* *cauchy-imp-bounded* **by** *fast*

have *AE* *x* *in* *M*. (λn . *diameter* {*s* *i* *x* | *i*. *n* \leq *i*}) \longrightarrow 0 **using** *Cauchy* *cauchy-iff-diameter-tends-to-zero-and-bounded* **by** *fast*

moreover **have** (λx . *diameter* {*s* *i* *x* | *i*. *n* \leq *i*}) \in *borel-measurable* *M* **for** *n*

using *bounded-range-s borel-measurable-diameter* **by** *measurable*
moreover have $AE\ x\ in\ M. \text{norm}(\text{diameter}\ \{s\ i\ x\ |\ i. n \leq i\}) \leq 4 * \text{norm}(f\ x)$
for n
proof –
{
 fix x **assume** $x: x \in \text{space } M$
 have $\text{diameter}\ \{s\ i\ x\ |\ i. n \leq i\} \leq 2 * \text{norm}(f\ x) + 2 * \text{norm}(f\ x)$
by (*intro diameter-le, blast, subst dist-norm[symmetric], intro dist-triangle3[THEN order-trans, of 0], intro add-mono*) (*auto intro: assms(5)[OF x]*)
 hence $\text{norm}(\text{diameter}\ \{s\ i\ x\ |\ i. n \leq i\}) \leq 4 * \text{norm}(f\ x)$ **using** *diameter-ge-0[OF bounded-subset[OF bounded-range-s], OF x, of {s i x | i. n ≤ i}]* **by** *force*
}
thus *?thesis* **by** *fast*
qed
ultimately have *diameter-tendsto-zero*: $(\lambda n. LINT\ x|M. \text{diameter}\ \{s\ i\ x\ |\ i. n \leq i\}) \longrightarrow 0$ **by** (*intro integral-dominated-convergence[OF borel-measurable-const[of 0] - integrable-4f, simplified]*) (*fast+*)

have *diameter-integrable*: *integrable* $M\ (\lambda x. \text{diameter}\ \{s\ i\ x\ |\ i. n \leq i\})$ **for** n
using *assms(1,5)* **by** (*intro integrable-bound-diameter[OF bounded-range-s integrable-2f], auto*)

have *dist-integrable*: *integrable* $M\ (\lambda x. \text{dist}(s\ i\ x)\ (s\ j\ x))$ **for** $i\ j$
using *assms(5) dist-triangle3[of s i - - 0, THEN order-trans, OF add-mono, of - 2 * norm(f -)]*
by (*intro Bochner-Integration.integrable-bound[OF integrable-4f]*) *fastforce+*

hence *dist-norm-integrable*: *integrable* $M\ (\lambda x. \text{norm}(s\ i\ x - s\ j\ x))$ **for** $i\ j$
unfolding *dist-norm* **by** *presburger*

have $\exists N. \forall i \geq N. \forall j \geq N. LINT\ x|M. \text{dist}(\text{cond-exp } M\ F\ (s\ i)\ x)\ (\text{cond-exp } M\ F\ (s\ j)\ x) < e$ **if** $e\text{-pos}: e > 0$ **for** e
proof –
 obtain N **where** $*$: $LINT\ x|M. \text{diameter}\ \{s\ i\ x\ |\ i. n \leq i\} < e$ **if** $n \geq N$ **for** n **using** *that order-tendsto-iff[THEN iffD1, OF diameter-tendsto-zero, unfolded eventually-sequentially] e-pos* **by** *presburger*
 {
 fix $i\ j\ x$ **assume** *asm*: $i \geq N\ j \geq N\ x \in \text{space } M$
 have $\text{case-prod dist } '(\{s\ i\ x\ |\ i. N \leq i\} \times \{s\ i\ x\ |\ i. N \leq i\}) = \text{case-prod } (\lambda i\ j. \text{dist}(s\ i\ x)\ (s\ j\ x))\ '(\{N..\} \times \{N..\})$ **by** *fast*
 hence $\text{diameter}\ \{s\ i\ x\ |\ i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist}(s\ i\ x)\ (s\ j\ x))$ **unfolding** *diameter-def* **by** *auto*
 moreover have $(\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist}(s\ i\ x)\ (s\ j\ x)) \geq \text{dist}(s\ i\ x)\ (s\ j\ x)$ **using** *asm bounded-imp-bdd-above[OF bounded-imp-dist-bounded, OF bounded-range-s]* **by** (*intro cSup-upper, auto*)
 ultimately have $\text{diameter}\ \{s\ i\ x\ |\ i. N \leq i\} \geq \text{dist}(s\ i\ x)\ (s\ j\ x)$ **by** *presburger*
 }
}

hence $LINT\ x|M. dist\ (s\ i\ x)\ (s\ j\ x) < e$ **if** $i \geq N\ j \geq N$ **for** $i\ j$ **using** *that* ***** **by** (*intro integral-mono*[*OF dist-integrable diameter-integrable, THEN order.strict-trans1*], *blast+*)

moreover have $LINT\ x|M. dist\ (cond-exp\ M\ F\ (s\ i)\ x)\ (cond-exp\ M\ F\ (s\ j)\ x) \leq LINT\ x|M. dist\ (s\ i\ x)\ (s\ j\ x)$ **for** $i\ j$

proof –

have $LINT\ x|M. dist\ (cond-exp\ M\ F\ (s\ i)\ x)\ (cond-exp\ M\ F\ (s\ j)\ x) = LINT\ x|M. norm\ (cond-exp\ M\ F\ (s\ i)\ x + -\ 1\ *_R\ cond-exp\ M\ F\ (s\ j)\ x)$ **unfolding** *dist-norm* **by** *simp*

also have $\dots = LINT\ x|M. norm\ (cond-exp\ M\ F\ (\lambda x. s\ i\ x - s\ j\ x)\ x)$ **using** *has-cond-exp-charact*(2)[*OF has-cond-exp-add*[*OF - has-cond-exp-scaleR-right, OF has-cond-exp-charact*(1,1), *OF has-cond-exp-simple*(1,1)[*OF assms*(2,3)]]], *THEN AE-symmetric, of i - 1 j*] **by** (*intro integral-cong-AE*) *force+*

also have $\dots \leq LINT\ x|M. cond-exp\ M\ F\ (\lambda x. norm\ (s\ i\ x - s\ j\ x))\ x$ **using** *cond-exp-contraction-simple*[*OF - fin-sup, of i j*] *integrable-cond-exp assms*(2) **by** (*intro integral-mono-AE, fast+*)

also have $\dots = LINT\ x|M. norm\ (s\ i\ x - s\ j\ x)$ **unfolding** *set-integral-space*(1)[*OF integrable-cond-exp, symmetric*] *set-integral-space*[*OF dist-norm-integrable, symmetric*] **by** (*intro has-cond-expD*(1)[*OF has-cond-exp-simple*[*OF - fin-sup-norm*], *symmetric*]) (*metis assms*(2) *simple-function-compose1 simple-function-diff, metis sets.top subalg subalgebra-def*)

finally show *?thesis unfolding dist-norm .*

qed

ultimately show *?thesis using order.strict-trans1 by meson*

qed

then obtain r **where** *strict-mono-r: strict-mono r* **and** *AE-Cauchy: AE x in M. Cauchy* ($\lambda i. cond-exp\ M\ F\ (s\ (r\ i))\ x$) **by** (*rule cauchy-L1-AE-cauchy-subseq*[*OF integrable-cond-exp*], *auto*)

hence *ae-lim-cond-exp: AE x in M. $(\lambda n. cond-exp\ M\ F\ (s\ (r\ n))\ x) \longrightarrow \lim$* ($\lambda n. cond-exp\ M\ F\ (s\ (r\ n))\ x$) **using** *Cauchy-convergent-iff convergent-LIMSEQ-iff* **by** *fastforce*

have *cond-exp-bounded: AE x in M. $norm\ (cond-exp\ M\ F\ (s\ (r\ n))\ x) \leq cond-exp\ M\ F\ (\lambda x. 2\ *\ norm\ (f\ x))\ x$* **for** n

proof –

have $AE\ x\ in\ M. norm\ (cond-exp\ M\ F\ (s\ (r\ n))\ x) \leq cond-exp\ M\ F\ (\lambda x. norm\ (s\ (r\ n)\ x))\ x$ **by** (*rule cond-exp-contraction-simple*[*OF assms*(2,3)])

moreover have $AE\ x\ in\ M. real-cond-exp\ M\ F\ (\lambda x. norm\ (s\ (r\ n)\ x))\ x \leq real-cond-exp\ M\ F\ (\lambda x. 2\ *\ norm\ (f\ x))\ x$ **using** *integrable-s integrable-2f assms*(5) **by** (*intro real-cond-exp-mono, auto*)

ultimately show *?thesis using cond-exp-real*[*OF integrable-norm, OF integrable-s, of r n*] *cond-exp-real*[*OF integrable-2f*] **by** *force*

qed

have *lim-integrable: integrable M $(\lambda x. \lim\ (\lambda i. cond-exp\ M\ F\ (s\ (r\ i))\ x))$* **by** (*intro integrable-dominated-convergence*[*OF - borel-measurable-cond-exp' integrable-cond-exp ae-lim-cond-exp cond-exp-bounded*], *simp*)

{
 fix A **assume** *A-in-sets-F: A ∈ sets F*

have $AE\ x\ in\ M. \text{norm}\ (\text{indicator}\ A\ x\ *_R\ \text{cond-exp}\ M\ F\ (s\ (r\ n))\ x) \leq \text{cond-exp}\ M\ F\ (\lambda x. 2 * \text{norm}\ (f\ x))\ x\ \text{for}\ n$
proof –
have $AE\ x\ in\ M. \text{norm}\ (\text{indicator}\ A\ x\ *_R\ \text{cond-exp}\ M\ F\ (s\ (r\ n))\ x) \leq \text{norm}\ (\text{cond-exp}\ M\ F\ (s\ (r\ n))\ x)\ \text{unfolding}\ \text{indicator-def}\ \text{by}\ \text{simp}$
thus $?thesis\ \text{using}\ \text{cond-exp-bounded}[of\ n]\ \text{by}\ \text{force}$
qed
hence $\text{lim-cond-exp-int}: (\lambda n. \text{LINT}\ x:A|M. \text{cond-exp}\ M\ F\ (s\ (r\ n))\ x) \longrightarrow \text{LINT}\ x:A|M. \text{lim}\ (\lambda n. \text{cond-exp}\ M\ F\ (s\ (r\ n))\ x)$
using $\text{ae-lim-cond-exp-measurable-from-subalg}[OF\ \text{subalg}\ \text{borel-measurable-indicator},\ OF\ A\text{-in-sets-}F]\ \text{cond-exp-bounded}$
unfolding $\text{set-lebesgue-integral-def}$
by $(\text{intro}\ \text{integral-dominated-convergence}[OF\ \text{borel-measurable-scaleR}\ \text{borel-measurable-scaleR}\ \text{integrable-cond-exp}])\ (\text{fastforce}\ \text{simp}\ \text{add:}\ \text{tendsto-scaleR})+$

have $AE\ x\ in\ M. \text{norm}\ (\text{indicator}\ A\ x\ *_R\ s\ (r\ n)\ x) \leq 2 * \text{norm}\ (f\ x)\ \text{for}\ n$
proof –
have $AE\ x\ in\ M. \text{norm}\ (\text{indicator}\ A\ x\ *_R\ s\ (r\ n)\ x) \leq \text{norm}\ (s\ (r\ n)\ x)\ \text{unfolding}\ \text{indicator-def}\ \text{by}\ \text{simp}$
thus $?thesis\ \text{using}\ \text{assms}(5)[of\ -\ r\ n]\ \text{by}\ \text{fastforce}$
qed
hence $\text{lim-s-int}: (\lambda n. \text{LINT}\ x:A|M. s\ (r\ n)\ x) \longrightarrow \text{LINT}\ x:A|M. f\ x$
using $\text{measurable-from-subalg}[OF\ \text{subalg}\ \text{borel-measurable-indicator},\ OF\ A\text{-in-sets-}F]\ \text{LIMSEQ-subseq-LIMSEQ}[OF\ \text{assms}(4)\ \text{strict-mono-r}]\ \text{assms}(5)$
unfolding $\text{set-lebesgue-integral-def}\ \text{comp-def}$
by $(\text{intro}\ \text{integral-dominated-convergence}[OF\ \text{borel-measurable-scaleR}\ \text{borel-measurable-scaleR}\ \text{integrable-2f}])\ (\text{fastforce}\ \text{simp}\ \text{add:}\ \text{tendsto-scaleR})+$

have $\text{LINT}\ x:A|M. \text{lim}\ (\lambda n. \text{cond-exp}\ M\ F\ (s\ (r\ n))\ x) = \text{lim}\ (\lambda n. \text{LINT}\ x:A|M. \text{cond-exp}\ M\ F\ (s\ (r\ n))\ x)\ \text{using}\ \text{limI}[OF\ \text{lim-cond-exp-int}]\ \text{by}\ \text{argo}$
also have $\dots = \text{lim}\ (\lambda n. \text{LINT}\ x:A|M. s\ (r\ n)\ x)\ \text{using}\ \text{has-cond-expD}(1)[OF\ \text{has-cond-exp-simple}[OF\ \text{assms}(2,3)]\ A\text{-in-sets-}F,\ \text{symmetric}]\ \text{by}\ \text{presburger}$
also have $\dots = \text{LINT}\ x:A|M. f\ x\ \text{using}\ \text{limI}[OF\ \text{lim-s-int}]\ \text{by}\ \text{argo}$
finally have $\text{LINT}\ x:A|M. \text{lim}\ (\lambda n. \text{cond-exp}\ M\ F\ (s\ (r\ n))\ x) = \text{LINT}\ x:A|M. f\ x.$
}
hence $\text{has-cond-exp}\ M\ F\ f\ (\lambda x. \text{lim}\ (\lambda i. \text{cond-exp}\ M\ F\ (s\ (r\ i))\ x))\ \text{using}\ \text{assms}(1)\ \text{lim-integrable}\ \text{by}\ (\text{intro}\ \text{has-cond-expI}',\ \text{auto})$
thus thesis using $AE\text{-Cauchy}\ \text{Cauchy-convergent}\ \text{strict-mono-r}\ \text{by}\ (\text{auto}\ \text{intro!}:\text{that})$
qed

lemma cond-exp-lim :

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology},\ \text{banach}\}$
assumes $[\text{measurable}]: \text{integrable}\ M\ f$
and $\bigwedge i. \text{simple-function}\ M\ (s\ i)$
and $\bigwedge i. \text{emeasure}\ M\ \{y \in \text{space}\ M. s\ i\ y \neq 0\} \neq \infty$
and $\bigwedge x. x \in \text{space}\ M \implies (\lambda i. s\ i\ x) \longrightarrow f\ x$
and $\bigwedge x\ i. x \in \text{space}\ M \implies \text{norm}\ (s\ i\ x) \leq 2 * \text{norm}\ (f\ x)$

obtains r **where** $AE\ x\ in\ M. (\lambda i. cond_exp\ M\ F\ (s\ (r\ i))\ x) \longrightarrow cond_exp\ M\ F\ f\ x\ strict_mono\ r$

proof –

obtain r **where** $AE\ x\ in\ M. cond_exp\ M\ F\ f\ x = \lim (\lambda i. cond_exp\ M\ F\ (s\ (r\ i))\ x) AE\ x\ in\ M. convergent (\lambda i. cond_exp\ M\ F\ (s\ (r\ i))\ x) strict_mono\ r$ **using** $has_cond_exp_charact(2)$ **by** $(auto\ intro: has_cond_exp_lim[OF\ assms])$
thus $?thesis$ **by** $(auto\ intro!: that[of\ r]\ simp: convergent-LIMSEQ-iff)$

qed

lemma has_cond_expI :

fixes $f :: 'a \Rightarrow 'b :: \{second_countable_topology, banach\}$

assumes $integrable\ M\ f$

shows $has_cond_exp\ M\ F\ f\ (cond_exp\ M\ F\ f)$

using $assms$

proof $(induction\ rule: integrable-induct')$

case $(base\ A\ c)$

show $?case$ **using** $has_cond_exp_indicator[OF\ base(1,2)]\ has_cond_exp_charact(1)$

by $blast$

next

case $(add\ u\ v)$

show $?case$ **using** $has_cond_exp_add[OF\ add(3,4)]\ has_cond_exp_charact(1)$ **by**

$blast$

next

case $(lim\ f\ s)$

show $?case$ **using** $has_cond_exp_lim[OF\ lim(1,3,4,5,6)]\ has_cond_exp_charact(1)$

by $meson$

qed

lemma $has_cond_exp_nested_subalg$:

fixes $f :: 'a \Rightarrow 'b :: \{second_countable_topology, banach\}$

assumes $subalgebra\ G\ F\ has_cond_exp\ M\ F\ f\ h\ has_cond_exp\ M\ G\ f\ h'$

shows $has_cond_exp\ M\ F\ h'\ h$

by $standard\ (metis\ assms\ has_cond_expD\ in-mono\ subalgebra-def)+$

lemma $cond_exp_nested_subalg$:

fixes $f :: 'a \Rightarrow 'b :: \{second_countable_topology, banach\}$

assumes $integrable\ M\ f\ subalgebra\ M\ G\ subalgebra\ G\ F$

shows $AE\ \xi\ in\ M. cond_exp\ M\ F\ f\ \xi = cond_exp\ M\ F\ (cond_exp\ M\ G\ f)\ \xi$

using $has_cond_expI\ assms\ sigma_finite_subalgebra_def$ **by** $(auto\ intro!: has_cond_exp_nested_subalg[THEN\ has_cond_exp_charact(2),\ THEN\ AE_symmetric]\ sigma_finite_subalgebra.has_cond_expI[OF\ sigma_finite_subalgebra.intro[OF\ assms(2)]]\ nested_subalg-is-sigma-finite)$

lemma $cond_exp_set_integral$:

fixes $f :: 'a \Rightarrow 'b :: \{second_countable_topology, banach\}$

assumes $integrable\ M\ f\ A \in sets\ F$

shows $(\int x \in A. f\ x\ \partial M) = (\int x \in A. cond_exp\ M\ F\ f\ x\ \partial M)$

using $has_cond_expD(1)[OF\ has_cond_expI,\ OF\ assms]$ **by** $argo$

lemma $cond_exp_add$:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes $\text{integrable } M f \text{ integrable } M g$
shows $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x + g x) x = \text{cond-exp } M F f x + \text{cond-exp } M F g x$
using $\text{has-cond-exp-add}[\text{OF has-cond-expI}(1,1), \text{OF assms}, \text{THEN has-cond-exp-charact}(2)]$
.

lemma *cond-exp-diff*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes $\text{integrable } M f \text{ integrable } M g$
shows $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x - g x) x = \text{cond-exp } M F f x - \text{cond-exp } M F g x$
using $\text{has-cond-exp-add}[\text{OF - has-cond-exp-scaleR-right}, \text{OF has-cond-expI}(1,1), \text{OF assms}, \text{THEN has-cond-exp-charact}(2), \text{of } -1]$ **by** *simp*

lemma *cond-exp-diff'*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes $\text{integrable } M f \text{ integrable } M g$
shows $AE x \text{ in } M. \text{cond-exp } M F (f - g) x = \text{cond-exp } M F f x - \text{cond-exp } M F g x$
unfolding *fun-diff-def* **using** *assms* **by** (*rule cond-exp-diff*)

lemma *cond-exp-contraction*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes $\text{integrable } M f$
shows $AE x \text{ in } M. \text{norm } (\text{cond-exp } M F f x) \leq \text{cond-exp } M F (\lambda x. \text{norm } (f x)) x$

proof –

obtain s **where** $s: \bigwedge i. \text{simple-function } M (s i) \bigwedge i. \text{emeasure } M \{y \in \text{space } M. s i y \neq 0\} \neq \infty \bigwedge x. x \in \text{space } M \implies (\lambda i. s i x) \longrightarrow f x \bigwedge i x. x \in \text{space } M \implies \text{norm } (s i x) \leq 2 * \text{norm } (f x)$
by (*blast intro: integrable-implies-simple-function-sequence* [*OF assms*])

obtain r **where** $r: AE x \text{ in } M. (\lambda i. \text{cond-exp } M F (s (r i)) x) \longrightarrow \text{cond-exp } M F f x \text{ strict-mono } r$ **using** *cond-exp-lim* [*OF assms* s] **by** *blast*

have *norm-s-r*: $\bigwedge i. \text{simple-function } M (\lambda x. \text{norm } (s (r i) x)) \bigwedge i. \text{emeasure } M \{y \in \text{space } M. \text{norm } (s (r i) y) \neq 0\} \neq \infty \bigwedge x. x \in \text{space } M \implies (\lambda i. \text{norm } (s (r i) x)) \longrightarrow \text{norm } (f x) \bigwedge i x. x \in \text{space } M \implies \text{norm } (\text{norm } (s (r i) x)) \leq 2 * \text{norm } (\text{norm } (f x))$

using s **by** (*auto intro: LIMSEQ-subseq-LIMSEQ* [*OF tendsto-norm* $r(2)$, *unfolded comp-def*] *simple-function-compose1*)

obtain r' **where** $r': AE x \text{ in } M. (\lambda i. (\text{cond-exp } M F (\lambda x. \text{norm } (s (r (r' i)) x))) x) \longrightarrow \text{cond-exp } M F (\lambda x. \text{norm } (f x)) x \text{ strict-mono } r'$ **using** *cond-exp-lim* [*OF integrable-norm norm-s-r*, *OF assms*] **by** *blast*

have $AE x \text{ in } M. \forall i. \text{norm } (\text{cond-exp } M F (s (r (r' i))) x) \leq \text{cond-exp } M F (\lambda x. \text{norm } (s (r (r' i)) x)) x$ **using** s **by** (*auto intro: cond-exp-contraction-simple simp*)

add: AE-all-countable)

moreover have $AE\ x\ in\ M.\ (\lambda i.\ norm\ (cond-exp\ M\ F\ (s\ (r\ (r'\ i))))\ x)) \longrightarrow$
 $norm\ (cond-exp\ M\ F\ f\ x)$ **using** $r\ LIMSEQ-subseq-LIMSEQ[OF\ tendsto-norm$
 $r'(2),\ unfolded\ comp-def]$ **by fast**
ultimately show *?thesis* **using** $LIMSEQ-le\ r'(1)$ **by fast**
qed

lemma *cond-exp-sum [intro, simp]:*

fixes $f :: 't \Rightarrow 'a \Rightarrow 'b :: \{second-countable-topology, banach\}$
assumes *[measurable]:* $\bigwedge i.\ integrable\ M\ (f\ i)$
shows $AE\ x\ in\ M.\ cond-exp\ M\ F\ (\lambda x.\ \sum_{i \in I}. f\ i\ x)\ x = (\sum_{i \in I}. cond-exp\ M\ F$
 $(f\ i)\ x)$
proof (*rule has-cond-exp-charact, intro has-cond-expI'*)
fix A **assume** *[measurable]:* $A \in sets\ F$
then have $A-meas\ [measurable]: A \in sets\ M$ **by** (*meson subsetD subalg subalgebra-def*)

have $(\int x \in A. (\sum_{i \in I}. f\ i\ x) \partial M) = (\int x. (\sum_{i \in I}. indicator\ A\ x\ *_R\ f\ i\ x) \partial M)$
unfolding *set-lebesgue-integral-def* **by** (*simp add: scaleR-sum-right*)
also have $\dots = (\sum_{i \in I}. (\int x. indicator\ A\ x\ *_R\ f\ i\ x\ \partial M))$ **using** *assms* **by** (*auto*
intro!: Bochner-Integration.integral-sum integrable-mult-indicator)
also have $\dots = (\sum_{i \in I}. (\int x. indicator\ A\ x\ *_R\ cond-exp\ M\ F\ (f\ i)\ x\ \partial M))$ **using**
 $cond-exp-set-integral[OF\ assms]$ **by** (*simp add: set-lebesgue-integral-def*)
also have $\dots = (\int x. (\sum_{i \in I}. indicator\ A\ x\ *_R\ cond-exp\ M\ F\ (f\ i)\ x) \partial M)$
using *assms* **by** (*auto intro!: Bochner-Integration.integral-sum[symmetric] integrable-mult-indicator*)
also have $\dots = (\int x \in A. (\sum_{i \in I}. cond-exp\ M\ F\ (f\ i)\ x) \partial M)$ **unfolding** *set-lebesgue-integral-def*
by (*simp add: scaleR-sum-right*)
finally show $(\int x \in A. (\sum_{i \in I}. f\ i\ x) \partial M) = (\int x \in A. (\sum_{i \in I}. cond-exp\ M\ F\ (f\ i)$
 $x) \partial M)$ **by auto**
qed (*auto simp add: assms integrable-cond-exp*)

1.5 Ordered Real Vectors

lemma *cond-exp-gr-c:*

fixes $f :: 'a \Rightarrow 'b :: \{second-countable-topology, banach, linorder-topology, ordered-real-vector\}$
assumes *integrable M f* $AE\ x\ in\ M.\ f\ x > c$
shows $AE\ x\ in\ M.\ cond-exp\ M\ F\ f\ x > c$
proof –
define X **where** $X = \{x \in space\ M.\ cond-exp\ M\ F\ f\ x \leq c\}$
have *[measurable]:* $X \in sets\ F$ **unfolding** $X-def$ **by** *measurable (metis sets.top subalg subalgebra-def)*
hence $X-in-M: X \in sets\ M$ **using** *sets-restr-to-subalg subalg subalgebra-def* **by blast**
have $emeasure\ M\ X = 0$
proof (*rule ccontr*)
assume $emeasure\ M\ X \neq 0$
have $emeasure\ (restr-to-subalg\ M\ F)\ X = emeasure\ M\ X$ **by** (*simp add: emea-*

sure-restr-to-subalg subalg)
hence *emeasure (restr-to-subalg M F) X > 0* **using** $\neg(\text{emeasure } M \ X = 0)$
gr-zeroI **by** *auto*
then obtain *A* **where** *A: A ∈ sets (restr-to-subalg M F) A ⊆ X* *emeasure*
(restr-to-subalg M F) A > 0 *emeasure (restr-to-subalg M F) A < ∞*
using *sigma-fin-subalg* **by** (*metis* *emeasure-notin-sets ennreal-0 infinity-ennreal-def*
le-less-linear *neg-top-trans* *not-gr-zero* *order-refl* *sigma-finite-measure.approx-PInf-emeasure-with-finite*)
hence [*simp*]: *A ∈ sets F* **using** *subalg sets-restr-to-subalg* **by** *blast*
hence [*simp*]: *A ∈ sets M* **using** *sets-restr-to-subalg subalg subalgebra-def* **by**
blast
have [*simp*]: *set-integrable M A (λx. c)* **using** *A subalg* **by** (*auto simp add:*
set-integrable-def *emeasure-restr-to-subalg*)
have [*simp*]: *set-integrable M A f* **unfolding** *set-integrable-def* **by** (*rule* *inte-*
grable-mult-indicator, auto simp add: assms(1))
have *AE x in M. indicator A x *_R c = indicator A x *_R f x*
proof (*rule integral-eq-mono-AE-eq-AE*)
show *LINT x|M. indicator A x *_R c = LINT x|M. indicator A x *_R f x*
proof (*simp only: set-lebesgue-integral-def[symmetric], rule antisym*)
show $(\int x \in A. c \ \partial M) \leq (\int x \in A. f \ x \ \partial M)$ **using** *assms(2)* **by** (*intro*
set-integral-mono-AE-banach) *auto*
have $(\int x \in A. f \ x \ \partial M) = (\int x \in A. \text{cond-exp } M \ F \ f \ x \ \partial M)$ **by** (*rule*
cond-exp-set-integral, auto simp add: <integrable M f>)
also have $\dots \leq (\int x \in A. c \ \partial M)$ **using** *A* **by** (*auto intro!: set-integral-mono-banach*
simp add: X-def)
finally show $(\int x \in A. f \ x \ \partial M) \leq (\int x \in A. c \ \partial M)$ **by** *simp*
qed
then have *measure M A *_R c = LINT x|M. indicator A x *_R f x* **using** *A*
by (*auto simp: set-lebesgue-integral-def* *emeasure-restr-to-subalg subalg*)
show *AE x in M. indicator A x *_R c ≤ indicator A x *_R f x* **using** *assms* **by**
(auto simp add: X-def indicator-def)
qed (*auto simp add: set-integrable-def[symmetric]*)
then have *AE x ∈ A in M. c = f x* **by** *auto*
then have *AE x ∈ A in M. False* **using** *assms(2)* **by** *auto*
have *A ∈ null-sets M* **unfolding** *ae-filter-def* **by** (*meson* *AE-iff-null-sets <A*
∈ sets M> <AE x ∈ A in M. False>)
then show *False* **using** *A(3)* **by** (*simp add: emeasure-restr-to-subalg null-setsD1*
subalg)
qed
then show *?thesis* **using** *AE-iff-null-sets[OF X-in-M]* **unfolding** *X-def* **by** *auto*
qed

lemma *cond-exp-less-c*:
fixes *f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-*
dered-real-vector}
assumes *integrable M f* *AE x in M. f x < c*
shows *AE x in M. cond-exp M F f x < c*
proof –
have *AE x in M. cond-exp M F f x = – cond-exp M F (λx. – f x) x* **using**
cond-exp-uminus[OF assms(1)] **by** *auto*

moreover have $AE\ x\ in\ M.\ cond_exp\ M\ F\ (\lambda x.\ -f\ x)\ x > -c$ **using** *assms* **by**
(intro cond-exp-gr-c) auto
ultimately show *?thesis* **by** *(force simp add: minus-less-iff)*
qed

lemma *cond-exp-mono-strict*:

fixes $f :: 'a \Rightarrow 'b :: \{second-countable-topology, banach, linorder-topology, ordered-real-vector\}$
assumes *integrable M f integrable M g AE x in M. f x < g x*
shows $AE\ x\ in\ M.\ cond_exp\ M\ F\ f\ x < cond_exp\ M\ F\ g\ x$
using *cond-exp-less-c[OF Bochner-Integration.integrable-diff, OF assms(1,2), of 0]*
cond-exp-diff[OF assms(1,2)] assms(3) **by** *auto*

lemma *cond-exp-ge-c*:

fixes $f :: 'a \Rightarrow 'b :: \{second-countable-topology, banach, linorder-topology, ordered-real-vector\}$
assumes [*measurable*]: *integrable M f*
and $AE\ x\ in\ M.\ f\ x \geq c$
shows $AE\ x\ in\ M.\ cond_exp\ M\ F\ f\ x \geq c$
proof –
let $?F = restr_to_subalg\ M\ F$
interpret *sigma-finite-measure restr-to-subalg M F* **using** *sigma-fin-subalg* **by**
auto
{
fix A **assume** *asm: A ∈ sets ?F 0 < measure ?F A*
have [*simp*]: *sets ?F = sets F measure ?F A = measure M A* **using** *asm* **by** (*auto simp add: measure-def sets-restr-to-subalg[OF subalg] emeasure-restr-to-subalg[OF subalg]*)
have $M-A: emeasure\ M\ A < \infty$ **using** *measure-zero-top asm* **by** (*force simp add: top.not-eq-extremum*)
hence $F-A: emeasure\ ?F\ A < \infty$ **using** *asm(1) emeasure-restr-to-subalg subalg*
by *fastforce*
have $set_lebesgue_integral\ M\ A\ (\lambda-. c) \leq set_lebesgue_integral\ M\ A\ f$ **using**
assms asm M-A subalg **by** (*intro set-integral-mono-AE-banach, auto simp add: set-integrable-def integrable-mult-indicator subalgebra-def sets-restr-to-subalg*)
also have $\dots = set_lebesgue_integral\ M\ A\ (cond_exp\ M\ F\ f)$ **using** *cond-exp-set-integral[OF assms(1)] asm* **by** *auto*
also have $\dots = set_lebesgue_integral\ ?F\ A\ (cond_exp\ M\ F\ f)$ **unfolding** *set-lebesgue-integral-def*
using *asm borel-measurable-cond-exp* **by** (*intro integral-subalgebra2[OF subalg, symmetric], simp*)
finally have $(1 / measure\ ?F\ A) *_R set_lebesgue_integral\ ?F\ A\ (cond_exp\ M\ F\ f)$
 $\in \{c..\}$ **using** *asm subalg M-A* **by** (*auto simp add: set-integral-const subalgebra-def intro!: pos-divideR-le-eq[THEN iffD1]*)
}
thus *?thesis* **using** *AE-restr-to-subalg[OF subalg] averaging-theorem[OF integrable-in-subalg closed-atLeast, OF subalg borel-measurable-cond-exp integrable-cond-exp]*
by *auto*
qed

lemma *cond-exp-le-c*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes *integrable* $M f$

and $AE\ x\ in\ M. f\ x \leq c$

shows $AE\ x\ in\ M. \text{cond-exp } M\ F\ f\ x \leq c$

proof –

have $AE\ x\ in\ M. \text{cond-exp } M\ F\ f\ x = -\ \text{cond-exp } M\ F\ (\lambda x. - f\ x)\ x$ **using** *cond-exp-uminus*[*OF assms(1)*] **by** *force*

moreover **have** $AE\ x\ in\ M. \text{cond-exp } M\ F\ (\lambda x. - f\ x)\ x \geq -\ c$ **using** *assms* **by** (*intro cond-exp-ge-c*) *auto*

ultimately show *?thesis* **by** (*force simp add: minus-le-iff*)

qed

lemma *cond-exp-mono*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes *integrable* $M f$ *integrable* $M g$ $AE\ x\ in\ M. f\ x \leq g\ x$

shows $AE\ x\ in\ M. \text{cond-exp } M\ F\ f\ x \leq \text{cond-exp } M\ F\ g\ x$

using *cond-exp-le-c*[*OF Bochner-Integration.integrable-diff, OF assms(1,2), of 0*]

cond-exp-diff[*OF assms(1,2)*] *assms(3)* **by** *auto*

lemma *cond-exp-min*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes *integrable* $M f$ *integrable* $M g$

shows $AE\ \xi\ in\ M. \text{cond-exp } M\ F\ (\lambda x. \min (f\ x)\ (g\ x))\ \xi \leq \min (\text{cond-exp } M\ F\ f\ \xi)\ (\text{cond-exp } M\ F\ g\ \xi)$

proof –

have $AE\ \xi\ in\ M. \text{cond-exp } M\ F\ (\lambda x. \min (f\ x)\ (g\ x))\ \xi \leq \text{cond-exp } M\ F\ f\ \xi$ **by** (*intro cond-exp-mono integrable-min assms, simp*)

moreover **have** $AE\ \xi\ in\ M. \text{cond-exp } M\ F\ (\lambda x. \min (f\ x)\ (g\ x))\ \xi \leq \text{cond-exp } M\ F\ g\ \xi$ **by** (*intro cond-exp-mono integrable-min assms, simp*)

ultimately show $AE\ \xi\ in\ M. \text{cond-exp } M\ F\ (\lambda x. \min (f\ x)\ (g\ x))\ \xi \leq \min (\text{cond-exp } M\ F\ f\ \xi)\ (\text{cond-exp } M\ F\ g\ \xi)$ **by** *fastforce*

qed

lemma *cond-exp-max*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes *integrable* $M f$ *integrable* $M g$

shows $AE\ \xi\ in\ M. \text{cond-exp } M\ F\ (\lambda x. \max (f\ x)\ (g\ x))\ \xi \geq \max (\text{cond-exp } M\ F\ f\ \xi)\ (\text{cond-exp } M\ F\ g\ \xi)$

proof –

have $AE\ \xi\ in\ M. \text{cond-exp } M\ F\ (\lambda x. \max (f\ x)\ (g\ x))\ \xi \geq \text{cond-exp } M\ F\ f\ \xi$ **by** (*intro cond-exp-mono integrable-max assms, simp*)

moreover **have** $AE\ \xi\ in\ M. \text{cond-exp } M\ F\ (\lambda x. \max (f\ x)\ (g\ x))\ \xi \geq \text{cond-exp } M\ F\ g\ \xi$

M F g ξ **by** (*intro cond-exp-mono integrable-max assms, simp*)
ultimately show *AE ξ in M. cond-exp M F (λx. max (f x) (g x)) ξ ≥ max*
(cond-exp M F f ξ) (cond-exp M F g ξ) **by** *fastforce*
qed

lemma *cond-exp-inf*:

fixes *f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, ordered-real-vector, lattice}*
assumes *integrable M f integrable M g*
shows *AE ξ in M. cond-exp M F (λx. inf (f x) (g x)) ξ ≤ inf (cond-exp M F f*
ξ) (cond-exp M F g ξ)
unfolding *inf-min* **using** *assms* **by** (*rule cond-exp-min*)

lemma *cond-exp-sup*:

fixes *f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, ordered-real-vector, lattice}*
assumes *integrable M f integrable M g*
shows *AE ξ in M. cond-exp M F (λx. sup (f x) (g x)) ξ ≥ sup (cond-exp M F f*
ξ) (cond-exp M F g ξ)
unfolding *sup-max* **using** *assms* **by** (*rule cond-exp-max*)

end

end

theory *Stochastic-Process*

imports *Filtration*

begin

1.6 Stochastic Process

locale *stochastic-process* = *sigma-finite-measure M* **for** *M* +

fixes *X :: 't :: {second-countable-topology, linorder-topology} ⇒ 'a ⇒ 'b :: {real-normed-vector, second-countable-topology}*

assumes *random-variable[measurable]: ∧i. X i ∈ borel-measurable M*

begin

definition *left-continuous* **where** *left-continuous* = (*AE ξ in M. ∀ i. continuous*
(at-left i) (λi. X i ξ))

definition *right-continuous* **where** *right-continuous* = (*AE ξ in M. ∀ i. continuous*
(at-right i) (λi. X i ξ))

lemma *compose*:

assumes *∧i. f i ∈ borel-measurable borel*

shows *stochastic-process M (λi ξ. (f i) (X i ξ))*

by (*unfold-locales, intro measurable-compose[OF random-variable assms]*)

lemma *norm: stochastic-process M (λi ξ. norm (X i ξ))* **by** (*auto intro: compose borel-measurable-norm*)

lemma *scaleR*:
assumes *stochastic-process* $M\ R$
shows *stochastic-process* $M\ (\lambda i\ \xi. (R\ i\ \xi) *_{\mathbb{R}} (X\ i\ \xi))$
by (*unfold-locales*) (*simp add: borel-measurable-scaleR random-variable assms stochastic-process.random-variable*)

lemma *scaleR-const-fun*:
assumes $f \in \text{borel-measurable } M$
shows *stochastic-process* $M\ (\lambda i\ \xi. f\ \xi *_{\mathbb{R}} (X\ i\ \xi))$
by (*unfold-locales, intro borel-measurable-scaleR assms random-variable*)

lemma *scaleR-const*: *stochastic-process* $M\ (\lambda i\ \xi. c *_{\mathbb{R}} (X\ i\ \xi))$ **by** (*auto intro: scaleR-const-fun borel-measurable-const*)

lemma *add*:
assumes *stochastic-process* $M\ Y$
shows *stochastic-process* $M\ (\lambda i\ \xi. X\ i\ \xi + Y\ i\ \xi)$
by (*unfold-locales*) (*simp add: borel-measurable-add random-variable assms stochastic-process.random-variable*)

lemma *diff*:
assumes *stochastic-process* $M\ Y$
shows *stochastic-process* $M\ (\lambda i\ \xi. X\ i\ \xi - Y\ i\ \xi)$
by (*unfold-locales*) (*simp add: borel-measurable-diff random-variable assms stochastic-process.random-variable*)

lemma *uminus*: *stochastic-process* $M\ (-X)$ **using** *scaleR-const*[*of -1*] **by** (*simp add: fun-Compl-def*)

end

1.7 Adapted Process

locale *adapted-process* = *filtered-sigma-finite-measure* $M\ F$ + *stochastic-process* $M\ X$ **for** M **and** $F :: 't :: \{\text{second-countable-topology, linorder-topology, order-bot}\} \Rightarrow$
- **and** $X :: 't \Rightarrow - \Rightarrow - :: \{\text{second-countable-topology, banach}\} +$
assumes *adapted*[*measurable*]: $\bigwedge i. X\ i \in \text{borel-measurable } (F\ i)$
begin

lemma *const-fun*:
assumes $f \in \text{borel-measurable } (F\ \text{bot})$
shows *adapted-process* $M\ F\ (\lambda -. f)$
using *assms* **by** (*unfold-locales*) (*blast intro: measurable-from-subalg subalgebra, metis borel-measurable-subalgebra bot.extremum sets-F-mono space-F*)

lemma *compose*:
assumes $\bigwedge i. f\ i \in \text{borel-measurable borel}$
shows *adapted-process* $M\ F\ (\lambda i\ \xi. (f\ i)\ (X\ i\ \xi))$
by (*unfold-locales, intro measurable-compose[OF random-variable assms], intro*

measurable-compose[*OF adapted assms*])

lemma *norm*: *adapted-process* $M F (\lambda i \xi. \text{norm } (X i \xi))$ **by** (*auto intro: compose borel-measurable-norm*)

lemma *scaleR*:

assumes *adapted-process* $M F R$

shows *adapted-process* $M F (\lambda i \xi. (R i \xi) *_{\mathbb{R}} (X i \xi))$

proof –

interpret R : *adapted-process* $M F R$ **by** (*rule assms*)

show ?thesis **by** (*unfold-locales*) (*auto simp add: borel-measurable-scaleR adapted random-variable assms R.random-variable R.adapted*)

qed

lemma *scaleR-const-fun*:

assumes $f \in \text{borel-measurable } (F \text{ bot})$

shows *adapted-process* $M F (\lambda i \xi. f \xi *_{\mathbb{R}} (X i \xi))$

using *assms* **by** (*fast intro: scaleR const-fun*)

lemma *scaleR-const*: *adapted-process* $M F (\lambda i \xi. c *_{\mathbb{R}} (X i \xi))$ **by** (*auto intro: scaleR-const-fun borel-measurable-const*)

lemma *add*:

assumes *adapted-process* $M F Y$

shows *adapted-process* $M F (\lambda i \xi. X i \xi + Y i \xi)$

proof –

interpret Y : *adapted-process* $M F Y$ **by** (*rule assms*)

show ?thesis **by** (*unfold-locales*) (*auto simp add: borel-measurable-add adapted random-variable Y.random-variable Y.adapted*)

qed

lemma *diff*:

assumes *adapted-process* $M F Y$

shows *adapted-process* $M F (\lambda i \xi. X i \xi - Y i \xi)$

proof –

interpret Y : *adapted-process* $M F Y$ **by** (*rule assms*)

show ?thesis **by** (*unfold-locales*) (*auto simp add: borel-measurable-diff adapted random-variable Y.random-variable Y.adapted*)

qed

lemma *uminus*: *adapted-process* $M F (-X)$ **using** *scaleR-const*[*of -1*] **by** (*simp add: fun-Compl-def*)

end

locale *adapted-process-order* = *adapted-process* $M F X$ **for** $M F$ **and** $X :: 't :: \{\text{second-countable-topology, linorder-topology, order-bot}\} \Rightarrow - \Rightarrow - :: \{\text{linorder-topology, ordered-real-vector}\}$

1.8 Discrete-Time Processes

locale *discrete-time-stochastic-process* = *stochastic-process* $M\ X$ **for** M **and** $X ::$
nat $\Rightarrow - \Rightarrow -$
locale *discrete-time-adapted-process* = *adapted-process* $M\ F\ X$ **for** $M\ F$ **and** $X ::$
nat $\Rightarrow - \Rightarrow -$
locale *discrete-time-adapted-process-order* = *adapted-process-order* $M\ F\ X$ **for** M
 F **and** $X ::$ *nat* $\Rightarrow - \Rightarrow -$

sublocale *discrete-time-adapted-process-order* \subseteq *discrete-time-adapted-process* **by**
(unfold-locale)
sublocale *discrete-time-adapted-process* \subseteq *discrete-time-stochastic-process* **by** (unfold-locale)
sublocale *discrete-time-adapted-process* \subseteq *nat-filtered-sigma-finite-measure* **by** (unfold-locale)

1.9 Predictable Processes

context *filtered-sigma-finite-measure*
begin

definition *predictable-sigma* :: $('t \times 'a)$ *measure* **where**
predictable-sigma = *sigma* ($UNIV \times$ *space* M) ($\{\{s <..t\} \times A \mid A\ s\ t.\ A \in F\ s \wedge$
 $s < t\} \cup \{\{bot\} \times A \mid A.\ A \in F\ bot\}\}$)

lemma *space-predictable-sigma[simp]*: *space* *predictable-sigma* = ($UNIV \times$ *space*
 M) **unfolding** *predictable-sigma-def* *space-measure-of-conv* **by** *blast*

lemma *sets-predictable-sigma[simp]*: *sets* *predictable-sigma* = *sigma-sets* ($UNIV \times$
space M) ($\{\{s <..t\} \times A \mid A\ s\ t.\ A \in F\ s \wedge s < t\} \cup \{\{bot\} \times A \mid A.\ A \in F\ bot\}\}$)
unfolding *predictable-sigma-def* *sets-measure-of-conv*
using *space-F* *sets.sets-into-space*
by (*fastforce* *intro!*: *if-P*)

definition *predictable* :: $('t \Rightarrow 'a \Rightarrow 'b :: \{second-countable-topology, banach\}) \Rightarrow$
bool **where**
predictable $X = (case-prod\ X \in borel-measurable\ (predictable-sigma))$

lemmas *predictableD* = *measurable-sets*[*OF* *predictable-def*[*THEN* *iffD1*], *unfolded*
space-predictable-sigma]

lemma (**in** *nat-filtered-sigma-finite-measure*) *predictable-sets-in-F*:
assumes $(\bigcup i.\ \{i\} \times A\ i) \in predictable-sigma$
shows $A\ (Suc\ i) \in F\ i$
 $A\ 0 \in F\ 0$
using *assms* **unfolding** *sets-predictable-sigma*
proof (*induction* $(\bigcup i.\ \{i\} \times A\ i)$ *arbitrary: A*)
case *Basic*
{
assume $\exists S.\ (\bigcup i.\ \{i\} \times A\ i) = \{bot\} \times S$
then obtain S **where** $S: (\bigcup i.\ \{i\} \times A\ i) = \{bot\} \times S$ **by** *blast*
hence $S \in F\ 0$ **using** *Basic* **by** (*fastforce* *simp* *add: times-eq-iff bot-nat-def*)
}

moreover have $A\ i = \{\}$ if $i \neq \text{bot}$ for i using that S by blast
 moreover have $A\ \text{bot} = S$ using S by blast
 ultimately have $A\ (\text{Suc } i) \in F\ i\ A\ 0 \in F\ 0$ for i unfolding bot-nat-def by
 (auto simp add: bot-nat-def)
 }
 note $\ast = \text{this}$
 {
 assume $\nexists S. (\bigcup i. \{i\} \times A\ i) = \{\text{bot}\} \times S$
 then obtain $s\ t\ B$ where $B: (\bigcup i. \{i\} \times A\ i) = \{s <.. t\} \times B\ B \in \text{sets } (F\ s)$
 $s < t$ using Basic by auto
 hence $A\ i = B$ if $i \in \{s <.. t\}$ for i using that by fast
 moreover have $A\ i = \{\}$ if $i \notin \{s <.. t\}$ for i using B that by fastforce
 ultimately have $A\ (\text{Suc } i) \in F\ i\ A\ 0 \in F\ 0$ for i unfolding bot-nat-def using
 $B\ \text{sets-F-mono}$ by (auto simp add: bot-nat-def) (metis less-Suc-eq-le sets.empty-sets
 subset-eq)
 }
 note $\ast\ast = \text{this}$
 show $A\ (\text{Suc } i) \in \text{sets } (F\ i)\ A\ 0 \in \text{sets } (F\ 0)$ using $\ast(1)[\text{of } i]\ \ast(2)\ \ast\ast(1)[\text{of } i]$
 $\ast\ast(2)$ by auto blast+
 next
 case Empty
 {
 case 1
 then show ?case using Empty by simp
 next
 case 2
 then show ?case using Empty by simp
 }
 next
 case (Compl a)
 have $a\text{-in}: a \subseteq \text{UNIV} \times \text{space } M$ using Compl(1) sets.sets-into-space sets-predictable-sigma
 space-predictable-sigma by metis
 hence $A\text{-in}: A\ i \subseteq \text{space } M$ for i using Compl(4) by blast
 have $a: a = \text{UNIV} \times \text{space } M - (\bigcup i. \{i\} \times A\ i)$ using $a\text{-in}$ Compl(4) by blast
 also have $\dots = (\bigcup j. \{j\} \times (\text{space } M - A\ j))$ by blast
 finally have $\ast: (\text{space } M - A\ (\text{Suc } i)) \in F\ i\ (\text{space } M - A\ 0) \in F\ 0$ using
 Compl(2,3) by auto
 {
 case 1
 then show ?case using $\ast\ A\text{-in}$ by (metis double-diff sets.compl-sets space-F
 subset-refl)
 next
 case 2
 then show ?case using $\ast\ A\text{-in}$ by (metis double-diff sets.compl-sets space-F
 subset-refl)
 }
 next
 case (Union a)
 have $a\text{-in}: a \subseteq \text{UNIV} \times \text{space } M$ for i using Union(1) sets.sets-into-space

sets-predictable-sigma space-predictable-sigma by metis
hence $A\text{-}i$: $A\ i \subseteq \text{space } M$ **for** i **using** $\text{Union}(4)$ **by** *blast*
have $\text{snd } x \in \text{snd } ' (a\ i \cap (\{fst\ x\} \times \text{space } M))$ **if** $x \in a\ i$ **for** $i\ x$ **using** *that*
a-in by fastforce
hence $a\text{-}i$: $a\ i = (\bigcup j. \{j\} \times (\text{snd } ' (a\ i \cap (\{j\} \times \text{space } M))))$ **for** i **by** *force*
have $A\text{-}i$: $A\ i = \text{snd } ' (\bigcup (\text{range } a) \cap (\{i\} \times \text{space } M))$ **for** i **unfolding**
 $\text{Union}(4)$ **using** $A\text{-}i$ **by** *force*
have $*$: $\text{snd } ' (a\ j \cap (\{Suc\ i\} \times \text{space } M)) \in F\ i\ \text{snd } ' (a\ j \cap (\{0\} \times \text{space } M))$
 $\in F\ 0$ **for** j **using** $\text{Union}(2,3)[OF\ a\text{-}i]$ **by** *auto*
 $\{$
 case 1
 have $(\bigcup j. \text{snd } ' (a\ j \cap (\{Suc\ i\} \times \text{space } M))) \in F\ i$ **using** $*$ **by** *fast*
 moreover **have** $(\bigcup j. \text{snd } ' (a\ j \cap (\{Suc\ i\} \times \text{space } M))) = \text{snd } ' (\bigcup (\text{range } a) \cap (\{Suc\ i\} \times \text{space } M))$ **by** *fast*
 ultimately show $?case$ **using** $A\text{-}i$ **by** *metis*
next
 case 2
 have $(\bigcup j. \text{snd } ' (a\ j \cap (\{0\} \times \text{space } M))) \in F\ 0$ **using** $*$ **by** *fast*
 moreover **have** $(\bigcup j. \text{snd } ' (a\ j \cap (\{0\} \times \text{space } M))) = \text{snd } ' (\bigcup (\text{range } a) \cap (\{0\} \times \text{space } M))$ **by** *fast*
 ultimately show $?case$ **using** $A\text{-}i$ **by** *metis*
 $\}$
qed

lemma (*in nat-filtered-sigma-finite-measure*) *predictable-discrete-time-process-measurable*:

assumes *predictable* X
shows $X\ i \in \text{borel-measurable } (F\ (i - 1))$
proof (*cases i*)
 case 0
 $\{$
 fix $S :: 'b\ \text{set}$ **assume** $\text{open-}S$: $\text{open } S$
 hence $\{0\} \times \text{space } M \in \text{predictable-sigma}$ **by** (*auto simp add: bot-nat-def space-F[symmetric, of bot]*)
 moreover **have** $\text{case-prod } X - ' S \cap (UNIV \times \text{space } M) \in \text{predictable-sigma}$
using $\text{open-}S$ **by** (*intro predictableD[OF assms], simp add: borel-open*)
 ultimately have $\text{case-prod } X - ' S \cap (\{0\} \times \text{space } M) \in \text{predictable-sigma}$
unfolding *sets-predictable-sigma* **using** $\text{space-}F\ \text{sets.sets-into-space}$
 by (*subst Times-Int-distrib1[of {0} UNIV space M, simplified], subst inf.commute[of - \times -], subst Int-assoc[symmetric], subst Int-range-binary*)
 (*intro sigma-sets-Inter binary-in-sigma-sets, fast*)
 moreover **have** $\text{case-prod } X - ' S \cap (\{0\} \times \text{space } M) = \{0\} \times (X\ 0 - ' S \cap \text{space } M)$ **by** (*auto simp add: le-Suc-eq*)
 moreover **have** $\dots = (\bigcup i. \{i\} \times (\text{if } i = 0 \text{ then } X\ 0 - ' S \cap \text{space } M \text{ else } \{\}))$
by (*auto split: if-splits*)
 ultimately have $(\bigcup i. \{i\} \times (\text{if } i = 0 \text{ then } X\ 0 - ' S \cap \text{space } M \text{ else } \{\})) \in \text{predictable-sigma}$ **by** *argo*
 then have $X\ 0 - ' S \cap \text{space } M \in \text{sets } (F\ 0)$ **using** *predictable-sets-in-F* [*of $\lambda i. \text{if } i = 0 \text{ then } X\ 0 - ' S \cap \text{space } M \text{ else } \{\}$*] **by** *presburger*
 $\}$

hence $X\ 0 \in \text{borel-measurable } (F\ 0)$ **by** (*fastforce simp add: bot-nat-def space-F intro!: borel-measurableI*)
 thus *?thesis* **using** 0 **by** *force*
next
 case ($Suc\ i$)
 {
 fix $S :: 'b\ set$ **assume** *open-S: open S*
 have $\{Suc\ i\} = \{i <.. Suc\ i\}$ **by** *fastforce*
 hence $\{Suc\ i\} \times \text{space } M \in \text{predictable-sigma}$ **unfolding** *space-F[symmetric, of i]* **by** (*auto intro!: sigma-sets.Basic*)
 moreover have $\text{case-prod } X - ' S \cap (UNIV \times \text{space } M) \in \text{predictable-sigma}$
using *open-S* **by** (*intro predictableD[OF assms], simp add: borel-open*)
 ultimately have $\text{case-prod } X - ' S \cap (\{Suc\ i\} \times \text{space } M) \in \text{predictable-sigma}$
unfolding *sets-predictable-sigma* **using** *space-F sets.sets-into-space*
by (*subst Times-Int-distrib1[of \{Suc i\} UNIV space M, simplified], subst inf.commute[of - \times -], subst Int-assoc[symmetric], subst Int-range-binary*)
 (*intro sigma-sets-Inter binary-in-sigma-sets, fast*) +
 moreover have $\text{case-prod } X - ' S \cap (\{Suc\ i\} \times \text{space } M) = \{Suc\ i\} \times (X\ (Suc\ i) - ' S \cap \text{space } M)$ **by** (*auto simp add: le-Suc-eq*)
 moreover have $\dots = (\bigcup j. \{j\} \times (\text{if } j = Suc\ i \text{ then } (X\ (Suc\ i) - ' S \cap \text{space } M) \text{ else } \{\})$ **by** (*auto split: if-splits*)
 ultimately have $(\bigcup j. \{j\} \times (\text{if } j = Suc\ i \text{ then } (X\ (Suc\ i) - ' S \cap \text{space } M) \text{ else } \{\})) \in \text{predictable-sigma}$ **by** *argo*
then have $X\ (Suc\ i) - ' S \cap \text{space } M \in \text{sets } (F\ i)$ **using** *predictable-sets-in-F[of \lambda j. if j = Suc i then (X (Suc i) - ' S \cap space M) else \{\}]* **by** *presburger*
 }
 hence $X\ (Suc\ i) \in \text{borel-measurable } (F\ i)$ **by** (*fastforce simp add: space-F intro!: borel-measurableI*)
then show *?thesis* **using** Suc **by** *force*
qed
end
end
theory *Martingale*
imports *Stochastic-Process Conditional-Expectation-Banach*
begin

1.10 Martingale

unbundle *lattice-syntax*

locale *martingale* = *adapted-process* +
assumes *integrable: $\bigwedge i. \text{integrable } M\ (X\ i)$*
and *martingale-property: $\bigwedge i\ j. i \leq j \implies AE\ \xi\ \text{in } M. X\ i\ \xi = \text{cond-exp } M\ (F\ i)\ (X\ j)\ \xi$*

lemma (*in filtered-sigma-finite-measure*) *martingale-const*[*intro*]:
assumes *integrable* $M\ f\ f \in \text{borel-measurable } (F\ \perp)$

shows *martingale* $M\ F\ (\lambda\cdot. f)$
using *assms cond-exp-F-meas*[*OF assms*(1), *THEN AE-symmetric*]
by (*unfold-locales*)
 (*simp add: borel-measurable-integrable,*
 metis bot.extremum measurable-from-subalg sets-F-mono space-F subalge-
bra-def, blast,
 metis (mono-tags, lifting) borel-measurable-subalgebra bot-least filtration.sets-F-mono
filtration-axioms space-F)

lemma (*in filtered-sigma-finite-measure*) *martingale-cond-exp*[*intro*]:
assumes *integrable* $M\ f$
shows *martingale* $M\ F\ (\lambda i. \text{cond-exp } M\ (F\ i)\ f)$
by (*unfold-locales,*
 auto simp add: subalgebra borel-measurable-cond-exp borel-measurable-cond-exp'
intro!: cond-exp-nested-subalg[*OF assms*],
 simp add: sets-F-mono space-F subalgebra-def)

1.11 Submartingale

locale *submartingale* = *adapted-process-order* +
assumes *integrable*: $\bigwedge i. \text{integrable } M\ (X\ i)$
and *submartingale-property*: $\bigwedge i\ j. i \leq j \implies AE\ \xi\ \text{in } M. X\ i\ \xi \leq \text{cond-exp } M\ (F\ i)\ (X\ j)\ \xi$

1.12 Supermartingale

locale *supermartingale* = *adapted-process-order* +
assumes *integrable*: $\bigwedge i. \text{integrable } M\ (X\ i)$
and *supermartingale-property*: $\bigwedge i\ j. i \leq j \implies AE\ \xi\ \text{in } M. X\ i\ \xi \geq \text{cond-exp } M\ (F\ i)\ (X\ j)\ \xi$

1.13 Martingale Stuff

locale *martingale-order* = *martingale* $M\ F\ X$ **for** $M\ F$ **and** $X :: - \Rightarrow - \Rightarrow - ::$
 $\{\text{linorder-topology, ordered-real-vector}\}$
begin

lemma *is-submartingale*: *submartingale* $M\ F\ X$ **using** *martingale-property* **by**
 (*unfold-locales*) (*force simp add: integrable*)+

lemma *is-supermartingale*: *supermartingale* $M\ F\ X$ **using** *martingale-property* **by**
 (*unfold-locales*) (*force simp add: integrable*)+

end

sublocale *martingale-order* \subseteq *martingale-is-submartingale*: *submartingale* **by** (*rule is-submartingale*)

sublocale *martingale-order* \subseteq *martingale-is-supermartingale*: *supermartingale* **by**
 (*rule is-supermartingale*)

```

locale submartingale-lattice = submartingale M F X for M F and X :: -  $\Rightarrow$  -  $\Rightarrow$ 
- :: {linorder-topology, lattice, ordered-real-vector}

locale supermartingale-lattice = supermartingale M F X for M F and X :: -  $\Rightarrow$  -
 $\Rightarrow$  - :: {linorder-topology, lattice, ordered-real-vector}

locale martingale-lattice = martingale M F X for M F and X :: -  $\Rightarrow$  -  $\Rightarrow$  - ::
{linorder-topology, lattice, ordered-real-vector}
begin

lemma is-submartingale: submartingale-lattice M F X using martingale-property
by (unfold-locales) (force simp add: integrable)+

lemma is-supermartingale: supermartingale-lattice M F X using martingale-property
by (unfold-locales) (force simp add: integrable)+

end

sublocale martingale-lattice  $\subseteq$  martingale-is-submartingale: submartingale-lattice
by (rule is-submartingale)

sublocale martingale-lattice  $\subseteq$  martingale-is-supermartingale: supermartingale-lattice
by (rule is-supermartingale)

context martingale
begin

lemma set-integral-eq:
  assumes A  $\in$  F i i  $\leq$  j
  shows set-lebesgue-integral M A (X i) = set-lebesgue-integral M A (X j)
proof -
  have  $\int x \in A. X i x \partial M = \int x \in A. \text{cond-exp } M (F i) (X j) x \partial M$  using
  martingale-property[OF assms(2)] borel-measurable-cond-exp' assms(1) subalgebra
  subalgebra-def by (intro set-lebesgue-integral-cong-AE[OF - random-variable]) fast-
  force+
  also have ... =  $\int x \in A. X j x \partial M$  using assms(1) by (auto simp: integrable
  intro: cond-exp-set-integral[symmetric])
  finally show ?thesis .
qed

lemma scaleR-const[intro]:
  shows martingale M F ( $\lambda i x. c *_R X i x$ )
proof -
  {
    fix i j :: 'b assume i  $\leq$  j
    hence AE x in M.  $c *_R X i x = \text{cond-exp } M (F i) (\lambda x. c *_R X j x) x$ 
    using cond-exp-scaleR-right[OF integrable, of i c, THEN AE-symmetric]
    martingale-property by force
  }

```

```

}
thus ?thesis by (unfold-locates) (auto simp add: borel-measurable-const-scaleR
adapted random-variable integrable)
qed

```

```

lemma uminus[intro]:
  shows martingale  $M\ F\ (-\ X)$ 
  using scaleR-const[of  $-1$ ] by (force intro: back-subst[of martingale  $M\ F$ ])

```

```

lemma add[intro]:
  assumes martingale  $M\ F\ Y$ 
  shows martingale  $M\ F\ (\lambda i\ \xi.\ X\ i\ \xi + Y\ i\ \xi)$ 
proof -
  interpret  $Y$ : martingale  $M\ F\ Y$  by (rule assms)
  {
    fix  $i\ j :: 'b$  assume asm:  $i \leq j$ 
    have  $AE\ \xi\ in\ M.\ X\ i\ \xi + Y\ i\ \xi = cond\_exp\ M\ (F\ i)\ (\lambda x.\ X\ j\ x + Y\ j\ x)\ \xi$ 
      using cond-exp-add[OF integrable martingale.integrable[OF assms], of  $i\ j\ j$ ,
      THEN AE-symmetric]
      martingale-property[OF asm] martingale.martingale-property[OF assms
asm] by force
  }
  thus ?thesis using assms
  by (unfold-locates) (auto simp add: borel-measurable-add random-variable adapted
integrable  $Y$ .adapted  $Y$ .random-variable martingale.integrable)
qed

```

```

lemma diff[intro]:
  assumes martingale  $M\ F\ Y$ 
  shows martingale  $M\ F\ (\lambda i\ x.\ X\ i\ x - Y\ i\ x)$ 
proof -
  interpret  $Y$ : martingale  $M\ F\ Y$  by (rule assms)
  {
    fix  $i\ j :: 'b$  assume asm:  $i \leq j$ 
    have  $AE\ \xi\ in\ M.\ X\ i\ \xi - Y\ i\ \xi = cond\_exp\ M\ (F\ i)\ (\lambda x.\ X\ j\ x - Y\ j\ x)\ \xi$ 
      using cond-exp-diff[OF integrable martingale.integrable[OF assms], of  $i\ j\ j$ ,
      THEN AE-symmetric, unfolded fun-diff-def]
      martingale-property[OF asm] martingale.martingale-property[OF assms
asm] by fastforce
  }
  thus ?thesis using assms by (unfold-locates) (auto simp add: borel-measurable-diff
random-variable adapted integrable  $Y$ .random-variable  $Y$ .adapted martingale.integrable)

```

qed

end

```

lemma (in adapted-process) martingale-of-set-integral-eq:
  assumes integrable:  $\bigwedge i.\ integrable\ M\ (X\ i)$ 

```

```

    and  $\bigwedge A \ i \ j. \ i \leq j \implies A \in F \ i \implies \text{set-lebesgue-integral } M \ A \ (X \ i) =$ 
    set-lebesgue-integral M A (X j)
    shows martingale M F X
  proof (unfold-locales)
    fix i j :: 't assume asm: i ≤ j
    interpret sigma-finite-measure restr-to-subalg M (F i) by (simp add: sigma-fin-subalg)
    {
      fix A assume A ∈ restr-to-subalg M (F i)
      hence *: A ∈ F i using sets-restr-to-subalg subalgebra by blast
      have set-lebesgue-integral (restr-to-subalg M (F i)) A (X i) = set-lebesgue-integral
      M A (X i) using * subalg by (auto simp: set-lebesgue-integral-def intro: inte-
      gral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator)
      also have ... = set-lebesgue-integral M A (cond-exp M (F i) (X j)) using *
      assms(2)[OF asm] cond-exp-set-integral[OF integrable] by auto
      finally have set-lebesgue-integral (restr-to-subalg M (F i)) A (X i) = set-lebesgue-integral
      (restr-to-subalg M (F i)) A (cond-exp M (F i) (X j)) using * subalg by (auto simp:
      set-lebesgue-integral-def intro!: integral-subalgebra2[symmetric] borel-measurable-scaleR
      borel-measurable-cond-exp borel-measurable-indicator)
    }
    hence AE ξ in restr-to-subalg M (F i). X i ξ = cond-exp M (F i) (X j) ξ by (intro
    density-unique, auto intro: integrable-in-subalg subalg borel-measurable-cond-exp in-
    tegrable)
    thus AE ξ in M. X i ξ = cond-exp M (F i) (X j) ξ using AE-restr-to-subalg[OF
    subalg] by blast
  qed (simp add: integrable)

```

```

lemma martingale-orderI:
  assumes submartingale M F X supermartingale M F X
  shows martingale-order M F X
proof -
  interpret submartingale M F X by (rule assms)
  interpret supermartingale M F X by (rule assms)
  show ?thesis using integrable submartingale-property supermartingale-property
  by (unfold-locales) (fast intro: antisym)+
qed

```

```

lemma martingale-iff: martingale M F X ⟷ submartingale M F X ∧ super-
martingale M F X
  using martingale-orderI martingale-order.is-submartingale martingale-order.is-supermartingale
  martingale-order-def by blast

```

1.14 Submartingale Stuff

```

context submartingale
begin

```

```

lemma set-integral-le:
  assumes A ∈ F i i ≤ j
  shows set-lebesgue-integral M A (X i) ≤ set-lebesgue-integral M A (X j)

```

```

unfolding cond-exp-set-integral[OF integrable assms(1), of j]
using submartingale-property[OF assms(2)]
by (simp only: set-lebesgue-integral-def, intro integral-mono-AE-banach, metis
assms(1) in-mono integrable integrable-mult-indicator subalgebra subalgebra-def,
metis assms(1) in-mono integrable-mult-indicator subalgebra subalgebra-def inte-
grable-cond-exp)
  (auto intro: scaleR-left-mono)

```

```

lemma cond-exp-diff-nonneg:
  assumes  $i \leq j$ 
  shows  $AE\ x\ in\ M. 0 \leq cond-exp\ M\ (F\ i)\ (\lambda \xi. X\ j\ \xi - X\ i\ \xi)\ x$ 
  using submartingale-property[OF assms] cond-exp-diff[OF integrable(1,1), of i j]
   $i]$  cond-exp-F-meas[OF integrable adapted, of i] by fastforce

```

```

lemma add[intro]:
  assumes submartingale  $M\ F\ Y$ 
  shows submartingale  $M\ F\ (\lambda i\ \xi. X\ i\ \xi + Y\ i\ \xi)$ 
proof -
  interpret  $Y$ : submartingale  $M\ F\ Y$  by (rule assms)
  {
    fix  $i\ j :: 'b$  assume  $asm: i \leq j$ 
    have  $AE\ \xi\ in\ M. X\ i\ \xi + Y\ i\ \xi \leq cond-exp\ M\ (F\ i)\ (\lambda x. X\ j\ x + Y\ j\ x)\ \xi$ 
    using cond-exp-add[OF integrable submartingale.integrable[OF assms], of i j]
  }
   $j]$ 
  submartingale-property[OF  $asm$ ] submartingale.submartingale-property[OF
assms  $asm$ ] add-mono[of  $X\ i - - Y\ i -$ ] by force
  }
  thus ?thesis using assms by (unfold-locale) (auto simp add: borel-measurable-add
random-variable adapted integrable  $Y$ .random-variable  $Y$ .adapted submartingale.integrable)

```

qed

```

lemma diff[intro]:
  assumes supermartingale  $M\ F\ Y$ 
  shows submartingale  $M\ F\ (\lambda i\ \xi. X\ i\ \xi - Y\ i\ \xi)$ 
proof -
  interpret  $Y$ : supermartingale  $M\ F\ Y$  by (rule assms)
  {
    fix  $i\ j :: 'b$  assume  $asm: i \leq j$ 
    have  $AE\ \xi\ in\ M. X\ i\ \xi - Y\ i\ \xi \leq cond-exp\ M\ (F\ i)\ (\lambda x. X\ j\ x - Y\ j\ x)\ \xi$ 
    using cond-exp-diff[OF integrable supermartingale.integrable[OF assms], of i
j j, unfolded fun-diff-def]
  }
  submartingale-property[OF  $asm$ ] supermartingale.supermartingale-property[OF
assms  $asm$ ] diff-mono[of  $X\ i - - Y\ i -$ ] by force
  }
  thus ?thesis using assms by (unfold-locale) (auto simp add: borel-measurable-diff
random-variable adapted integrable  $Y$ .random-variable  $Y$ .adapted supermartingale.integrable)

```

qed

lemma *scaleR-nonneg*:
assumes $c \geq 0$
shows *submartingale* $M F (\lambda i \xi. c *_{\mathbb{R}} X i \xi)$
proof
{
 fix $i j :: 'b$ **assume** $asm: i \leq j$
 show $AE \xi \text{ in } M. c *_{\mathbb{R}} X i \xi \leq \text{cond-exp } M (F i) (\lambda \xi. c *_{\mathbb{R}} X j \xi) \xi$
 using *cond-exp-scaleR-right*[*OF integrable, of i c j*] *submartingale-property*[*OF asm*] **by** (*auto intro!*: *scaleR-left-mono*[*OF - assms*])
}
qed (*auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable random-variable adapted borel-measurable-const-scaleR*)

lemma *scaleR-nonpos*:
assumes $c \leq 0$
shows *supermartingale* $M F (\lambda i \xi. c *_{\mathbb{R}} X i \xi)$
proof
{
 fix $i j :: 'b$ **assume** $asm: i \leq j$
 show $AE \xi \text{ in } M. c *_{\mathbb{R}} X i \xi \geq \text{cond-exp } M (F i) (\lambda \xi. c *_{\mathbb{R}} X j \xi) \xi$
 using *cond-exp-scaleR-right*[*OF integrable, of i c j*] *submartingale-property*[*OF asm*] **by** (*auto intro!*: *scaleR-left-mono-neg*[*OF - assms*])
}
qed (*auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable random-variable adapted borel-measurable-const-scaleR*)

lemma *uminus*[*intro*]:
shows *supermartingale* $M F (- X)$
unfolding *fun-Compl-def* **using** *scaleR-nonpos*[*of -1*] **by** *simp*

lemma *max*:
assumes *submartingale* $M F Y$
shows *submartingale* $M F (\lambda i \xi. \max (X i \xi) (Y i \xi))$
proof (*unfold-locales*)
 interpret $Y: \text{submartingale } M F Y$ **by** (*rule assms*)
 {
 fix $i j :: 'b$ **assume** $asm: i \leq j$
 have $AE \xi \text{ in } M. \max (X i \xi) (Y i \xi) \leq \max (\text{cond-exp } M (F i) (X j) \xi) (\text{cond-exp } M (F i) (Y j) \xi)$ **using** *submartingale-property* $Y.\text{submartingale-property } asm$ **unfolding** *max-def* **by** *fastforce*
 thus $AE \xi \text{ in } M. \max (X i \xi) (Y i \xi) \leq \text{cond-exp } M (F i) (\lambda \xi. \max (X j \xi) (Y j \xi)) \xi$ **using** *cond-exp-max*[*OF integrable Y.integrable, of i j j*] *order.trans* **by** *fast*
 }
 show $\bigwedge i. (\lambda \xi. \max (X i \xi) (Y i \xi)) \in \text{borel-measurable } M \bigwedge i. (\lambda \xi. \max (X i \xi) (Y i \xi)) \in \text{borel-measurable } (F i) \bigwedge i. \text{integrable } M (\lambda \xi. \max (X i \xi) (Y i \xi))$ **by** (*force intro: Y.integrable integrable assms*)
qed

```

lemma max-0:
  shows submartingale  $M F (\lambda i \xi. \max 0 (X i \xi))$ 
proof -
  interpret zero: submartingale  $M F \lambda -. 0$  by (intro martingale-order.is-submartingale,
unfold-locales, auto)
  show ?thesis by (intro zero.max submartingale-axioms)
qed

end

lemma (in submartingale-lattice) sup:
  assumes submartingale-lattice  $M F Y$ 
  shows submartingale-lattice  $M F (\lambda i \xi. \sup (X i \xi) (Y i \xi))$ 
  using submartingale-lattice.intro submartingale.max [OF submartingale-axioms
assms [THEN submartingale-lattice.axioms]] unfolding sup-max [symmetric] .

lemma (in adapted-process-order) submartingale-of-cond-exp-diff-nonneg:
  assumes integrable:  $\bigwedge i. \text{integrable } M (X i)$ 
  and diff-nonneg:  $\bigwedge i j. i \leq j \implies \text{AE } x \text{ in } M. 0 \leq \text{cond-exp } M (F i) (\lambda \xi. X j$ 
 $\xi - X i \xi) x$ 
  shows submartingale  $M F X$ 
proof (unfold-locales)
  {
    fix  $i j :: 't$  assume asm:  $i \leq j$ 
    show  $\text{AE } \xi \text{ in } M. X i \xi \leq \text{cond-exp } M (F i) (X j) \xi$ 
    using diff-nonneg [OF asm] cond-exp-diff [OF integrable (1,1), of i j i] cond-exp-F-meas [OF
integrable adapted, of i] by fastforce
  }
qed (intro integrable)

lemma (in adapted-process-order) submartingale-of-set-integral-le:
  assumes integrable:  $\bigwedge i. \text{integrable } M (X i)$ 
  and  $\bigwedge A i j. i \leq j \implies A \in F i \implies \text{set-lebesgue-integral } M A (X i) \leq$ 
 $\text{set-lebesgue-integral } M A (X j)$ 
  shows submartingale  $M F X$ 
proof (unfold-locales)
  {
    fix  $i j :: 't$  assume asm:  $i \leq j$ 
    interpret sigma-finite-measure restr-to-subalg  $M (F i)$  by (simp add: sigma-fin-subalg)
    {
      fix  $A$  assume  $A \in \text{restr-to-subalg } M (F i)$ 
      hence *:  $A \in F i$  using sets-restr-to-subalg subalgebra by blast
      have set-lebesgue-integral (restr-to-subalg  $M (F i)$ )  $A (X i) = \text{set-lebesgue-integral}$ 
 $M A (X i)$  using * subalg by (auto simp: set-lebesgue-integral-def intro: inte-
gral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator)
      also have ...  $\leq \text{set-lebesgue-integral } M A (\text{cond-exp } M (F i) (X j))$  using *
assms (2) [OF asm] cond-exp-set-integral [OF integrable] by auto
      also have ...  $= \text{set-lebesgue-integral } (\text{restr-to-subalg } M (F i)) A (\text{cond-exp}$ 

```

$M (F i) (X j))$ **using** * *subalg* **by** (*auto simp: set-lebesgue-integral-def intro!: integral-subalgebra2[symmetric] borel-measurable-scaleR borel-measurable-cond-exp borel-measurable-indicator*)
finally have $0 \leq \text{set-lebesgue-integral } (\text{restr-to-subalg } M (F i)) A (\lambda \xi. \text{cond-exp } M (F i) (X j) \xi - X i \xi)$ **using** * *subalg* **by** (*subst set-integral-diff, auto simp add: set-integrable-def sets-restr-to-subalg intro!: integrable adapted integrable-in-subalg borel-measurable-scaleR borel-measurable-indicator borel-measurable-cond-exp integrable-mult-indicator*)
}
hence $AE \xi \text{ in } \text{restr-to-subalg } M (F i). 0 \leq \text{cond-exp } M (F i) (X j) \xi - X i \xi$ **by** (*intro density-nonneg integrable-in-subalg subalg borel-measurable-diff borel-measurable-cond-exp adapted Bochner-Integration.integrable-diff integrable-cond-exp integrable*)
thus $AE \xi \text{ in } M. X i \xi \leq \text{cond-exp } M (F i) (X j) \xi$ **using** *AE-restr-to-subalg[OF subalg]* **by** *simp*
}
qed (*intro integrable*)

1.15 Supermartingale Stuff

context *supermartingale*
begin

lemma *set-integral-ge*:
assumes $A \in F i i \leq j$
shows $\text{set-lebesgue-integral } M A (X i) \geq \text{set-lebesgue-integral } M A (X j)$
unfolding *cond-exp-set-integral[OF integrable assms(1), of j]*
using *supermartingale-property[OF assms(2)]*
by (*simp only: set-lebesgue-integral-def, intro integral-mono-AE-banach, metis assms(1) in-mono integrable-mult-indicator subalgebra subalgebra-def integrable-cond-exp, metis assms(1) in-mono integrable integrable-mult-indicator subalgebra subalgebra-def*)
(auto intro: scaleR-left-mono)

lemma *cond-exp-diff-nonneg*:
assumes $i \leq j$
shows $AE x \text{ in } M. 0 \leq \text{cond-exp } M (F i) (\lambda \xi. X i \xi - X j \xi) x$
using *supermartingale-property[OF assms] cond-exp-diff[OF integrable(1,1), of i i j] cond-exp-F-meas[OF integrable adapted, of i]* **by** *fastforce*

lemma *add[intro]*:
assumes *supermartingale M F Y*
shows *supermartingale M F* $(\lambda i \xi. X i \xi + Y i \xi)$
proof –
interpret *Y: supermartingale M F Y* **by** (*rule assms*)
{
fix $i j :: 'b$ **assume** *asm: i ≤ j*
have $AE \xi \text{ in } M. X i \xi + Y i \xi \geq \text{cond-exp } M (F i) (\lambda x. X j x + Y j x) \xi$
using *cond-exp-add[OF integrable supermartingale.integrable[OF assms], of i j j]*

```

      supermartingale-property[OF asm] supermartingale.supermartingale-property[OF
assms asm] add-mono[of - X i - - Y i -] by force
    }
    thus ?thesis using assms by (unfold-locale) (auto simp add: borel-measurable-add
random-variable adapted integrable Y.random-variable Y.adapted supermartingale.integrable)

qed

```

```

lemma diff[intro]:
  assumes submartingale M F Y
  shows supermartingale M F ( $\lambda i \xi. X i \xi - Y i \xi$ )
proof -
  interpret Y: submartingale M F Y by (rule assms)
  {
    fix i j :: 'b assume asm:  $i \leq j$ 
    have AE  $\xi$  in M.  $X i \xi - Y i \xi \geq \text{cond-exp } M (F i) (\lambda x. X j x - Y j x) \xi$ 
      using cond-exp-diff[OF integrable submartingale.integrable[OF assms], of i j
j, unfolded fun-diff-def]
      supermartingale-property[OF asm] submartingale.submartingale-property[OF
assms asm] diff-mono[of - X i - Y i -] by force
    }
    thus ?thesis using assms by (unfold-locale) (auto simp add: borel-measurable-diff
random-variable adapted integrable Y.random-variable Y.adapted submartingale.integrable)

qed

```

```

lemma scaleR-nonneg:
  assumes  $c \geq 0$ 
  shows supermartingale M F ( $\lambda i \xi. c *_R X i \xi$ )
proof
  {
    fix i j :: 'b assume asm:  $i \leq j$ 
    show AE  $\xi$  in M.  $c *_R X i \xi \geq \text{cond-exp } M (F i) (\lambda \xi. c *_R X j \xi) \xi$ 
      using cond-exp-scaleR-right[OF integrable, of i c j] supermartingale-property[OF
asm]
      by (auto intro!: scaleR-left-mono[OF - assms])
    }
  qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)

```

```

lemma scaleR-nonpos:
  assumes  $c \leq 0$ 
  shows submartingale M F ( $\lambda i \xi. c *_R X i \xi$ )
proof
  {
    fix i j :: 'b assume asm:  $i \leq j$ 
    show AE  $\xi$  in M.  $c *_R X i \xi \leq \text{cond-exp } M (F i) (\lambda \xi. c *_R X j \xi) \xi$ 
      using cond-exp-scaleR-right[OF integrable, of i c j] supermartingale-property[OF
asm]

```

```

    by (auto intro!: scaleR-left-mono-neg[OF - assms])
  }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)

lemma uminus[intro]:
  shows submartingale M F (− X)
  unfolding fun-Compl-def using scaleR-nonpos[of −1] by simp

lemma min:
  assumes supermartingale M F Y
  shows supermartingale M F (λi ξ. min (X i ξ) (Y i ξ))
proof (unfold-locale)
  interpret Y: supermartingale M F Y by (rule assms)
  {
    fix i j :: 'b assume asm: i ≤ j
    have AE ξ in M. min (X i ξ) (Y i ξ) ≥ min (cond-exp M (F i) (X j) ξ) (cond-exp
M (F i) (Y j) ξ) using supermartingale-property Y.supermartingale-property asm
    unfolding min-def by fastforce
    thus AE ξ in M. min (X i ξ) (Y i ξ) ≥ cond-exp M (F i) (λξ. min (X j ξ) (Y
j ξ)) ξ using cond-exp-min[OF integrable Y.integrable, of i j j] order.trans by fast
  }
  show ∧i. (λξ. min (X i ξ) (Y i ξ)) ∈ borel-measurable M ∧i. (λξ. min (X i ξ)
(Y i ξ)) ∈ borel-measurable (F i) ∧i. integrable M (λξ. min (X i ξ) (Y i ξ)) by
(force intro: Y.integrable integrable assms)+
qed

lemma min-0:
  shows supermartingale M F (λi ξ. min 0 (X i ξ))
proof −
  interpret zero: supermartingale M F λ-. 0 by (intro martingale-order.is-supermartingale,
unfold-locale, auto)
  show ?thesis by (intro zero.min supermartingale-axioms)
qed

end

lemma (in supermartingale-lattice) inf:
  assumes supermartingale-lattice M F Y
  shows supermartingale-lattice M F (λi ξ. inf (X i ξ) (Y i ξ))
  using supermartingale-lattice.intro supermartingale.min[OF supermartingale-axioms
assms[THEN supermartingale-lattice.axioms]] unfolding inf-min[symmetric] .

lemma (in adapted-process-order) supermartingale-of-cond-exp-diff-nonneg:
  assumes integrable: ∧i. integrable M (X i)
  and diff-nonneg: ∧i j. i ≤ j ⟹ AE x in M. 0 ≤ cond-exp M (F i) (λξ. X i
ξ − X j ξ) x
  shows supermartingale M F X
proof

```

```

{
  fix i j :: 't assume asm: i ≤ j
  show AE ξ in M. X i ξ ≥ cond-exp M (F i) (X j) ξ
  using diff-nonneg[OF asm] cond-exp-diff[OF integrable(1,1), of i i j] cond-exp-F-meas[OF
integrable adapted, of i] by fastforce
}
qed (intro integrable)

```

```

lemma (in adapted-process-order) supermartingale-of-set-integral-ge:
  assumes integrable:  $\bigwedge i. \text{integrable } M (X i)$ 
    and  $\bigwedge A i j. i \leq j \implies A \in F i \implies \text{set-lebesgue-integral } M A (X j) \leq$ 
 $\text{set-lebesgue-integral } M A (X i)$ 
  shows supermartingale M F X
proof -
  interpret uminus-X: adapted-process-order M F - X by (intro adapted-process-order.intro
uminus)
  note * = set-integral-uminus[unfolded set-integrable-def, OF integrable-mult-indicator[OF
- integrable]]
  have supermartingale M F (-(- X)) using ord-eq-le-trans[OF * ord-le-eq-trans[OF
le-imp-neg-le[OF assms(2)] *[symmetric]]] subalg
  by (intro submartingale.uminus uminus-X.submartingale-of-set-integral-le) (auto
simp add: subalgebra-def integrable fun-Compl-def, blast)
  thus ?thesis unfolding fun-Compl-def by simp
qed

```

2 Discrete Time Martingales

```

locale discrete-time-martingale = martingale M F X for M F and X :: nat  $\Rightarrow$  -
 $\Rightarrow$  -
locale discrete-time-submartingale = submartingale M F X for M F and X :: nat
 $\Rightarrow$  -  $\Rightarrow$  -
locale discrete-time-supermartingale = supermartingale M F X for M F and X
:: nat  $\Rightarrow$  -  $\Rightarrow$  -

```

```

sublocale discrete-time-martingale  $\subseteq$  discrete-time-adapted-process by (unfold-locales)
sublocale discrete-time-submartingale  $\subseteq$  discrete-time-adapted-process by (unfold-locales)
sublocale discrete-time-supermartingale  $\subseteq$  discrete-time-adapted-process by (unfold-locales)

```

3 Discrete Time Martingales

```

lemma (in discrete-time-martingale) predictable-eq-bot:
  assumes predictable X
  shows AE ξ in M. X i ξ = X  $\perp$  ξ
proof (induction i)
  case 0
  then show ?case by (simp add: bot-nat-def)
next
  case (Suc i)

```

thus ?case **using** predictable-discrete-time-process-measurable[OF assms, of Suc
i]
 martingale-property[OF le-SucI, of *i*]
 cond-exp-F-meas[OF integrable, of Suc *i i*] Suc **by** fastforce
qed

lemma (in discrete-time-adapted-process) martingale-of-set-integral-eq-Suc:
 assumes integrable: $\bigwedge i. \text{integrable } M \ (X \ i)$
 and $\bigwedge A \ i. A \in F \ i \implies \text{set-lebesgue-integral } M \ A \ (X \ i) = \text{set-lebesgue-integral}$
 $M \ A \ (X \ (\text{Suc } i))$
 shows discrete-time-martingale $M \ F \ X$
proof (intro discrete-time-martingale.intro martingale-of-set-integral-eq)
 fix *i j A* **assume** asm: $i \leq j \ A \in \text{sets } (F \ i)$
 show $\text{set-lebesgue-integral } M \ A \ (X \ i) = \text{set-lebesgue-integral } M \ A \ (X \ j)$ **using**
 asm
proof (induction $j - i$ arbitrary: *i j*)
 case 0
 then show ?case **using** asm **by** simp
 next
 case (Suc *n*)
 hence *: $n = j - \text{Suc } i$ **by** linarith
 have $\text{Suc } i \leq j$ **using** Suc(2,3) **by** linarith
 thus ?case **using** sets-F-mono[OF le-SucI] Suc(4) Suc(1)[OF *] **by** (auto
 intro: assms(2)[THEN trans])
qed
qed (simp add: integrable)

lemma (in discrete-time-adapted-process) martingale-nat:
 assumes integrable: $\bigwedge i. \text{integrable } M \ (X \ i)$
 and $\bigwedge i. AE \ \xi \text{ in } M. X \ i \ \xi = \text{cond-exp } M \ (F \ i) \ (X \ (\text{Suc } i)) \ \xi$
 shows discrete-time-martingale $M \ F \ X$
proof (unfold-locales)
 fix *i j :: nat* **assume** asm: $i \leq j$
 show $AE \ \xi \text{ in } M. X \ i \ \xi = \text{cond-exp } M \ (F \ i) \ (X \ j) \ \xi$ **using** asm
proof (induction $j - i$ arbitrary: *i j*)
 case 0
 hence $j = i$ **by** simp
 thus ?case **using** cond-exp-F-meas[OF integrable adapted, THEN AE-symmetric]
by presburger
 next
 case (Suc *n*)
 have $j: j = \text{Suc } (n + i)$ **using** Suc **by** linarith
 have $n: n = n + i - i$ **using** Suc **by** linarith
 have *: $AE \ \xi \text{ in } M. \text{cond-exp } M \ (F \ (n + i)) \ (X \ j) \ \xi = X \ (n + i) \ \xi$ **unfolding**
 j **using** assms(2)[THEN AE-symmetric] **by** blast
 have $AE \ \xi \text{ in } M. \text{cond-exp } M \ (F \ i) \ (X \ j) \ \xi = \text{cond-exp } M \ (F \ i) \ (\text{cond-exp } M$
 $(F \ (n + i)) \ (X \ j)) \ \xi$ **by** (intro cond-exp-nested-subalg integrable subalg, simp add:
 subalgebra-def space-F sets-F-mono)
 hence $AE \ \xi \text{ in } M. \text{cond-exp } M \ (F \ i) \ (X \ j) \ \xi = \text{cond-exp } M \ (F \ i) \ (X \ (n + i))$

ξ **using** *cond-exp-cong-AE*[*OF integrable-cond-exp integrable **] **by** *force*
 thus ?case **using** *Suc(1)*[*OF n*] **by** *fastforce*
qed
qed (*simp add: integrable*)

lemma (*in discrete-time-adapted-process*) *martingale-of-cond-exp-diff-Suc-eq-0*:
 assumes *integrable*: $\bigwedge i. \text{integrable } M \ (X \ i)$
 and $\bigwedge i. AE \ \xi \text{ in } M. 0 = \text{cond-exp } M \ (F \ i) \ (\lambda \xi. X \ (Suc \ i) \ \xi - X \ i \ \xi) \ \xi$
 shows *discrete-time-martingale* *M F X*
proof (*intro martingale-nat integrable*)
 fix *i*
 show $AE \ \xi \text{ in } M. X \ i \ \xi = \text{cond-exp } M \ (F \ i) \ (X \ (Suc \ i)) \ \xi$ **using** *cond-exp-diff*[*OF integrable(1,1), of i Suc i i*] *cond-exp-F-meas*[*OF integrable adapted, of i*] *assms(2)*[*of i*] **by** *fastforce*
qed

4 Discrete Time Submartingales

lemma (*in discrete-time-submartingale*) *predictable-ge-bot*:
 assumes *predictable* *X*
 shows $AE \ \xi \text{ in } M. X \ i \ \xi \geq X \ \perp \ \xi$
proof (*induction i*)
 case 0
 then show ?case **by** (*simp add: bot-nat-def*)
next
 case (*Suc i*)
 thus ?case **using** *predictable-discrete-time-process-measurable*[*OF assms, of Suc i*]
 $\text{submartingale-property}$ [*OF le-SucI, of i*]
 cond-exp-F-meas [*OF integrable, of Suc i i*] *Suc* **by** *fastforce*
qed

lemma (*in discrete-time-adapted-process-order*) *submartingale-of-set-integral-le-Suc*:
 assumes *integrable*: $\bigwedge i. \text{integrable } M \ (X \ i)$
 and $\bigwedge A \ i. A \in F \ i \implies \text{set-lebesgue-integral } M \ A \ (X \ i) \leq \text{set-lebesgue-integral } M \ A \ (X \ (Suc \ i))$
 shows *discrete-time-submartingale* *M F X*
proof (*intro discrete-time-submartingale.intro submartingale-of-set-integral-le*)
 fix *i j A* **assume** *asm*: $i \leq j \wedge A \in \text{sets } (F \ i)$
 show $\text{set-lebesgue-integral } M \ A \ (X \ i) \leq \text{set-lebesgue-integral } M \ A \ (X \ j)$ **using** *asm*
proof (*induction j - i arbitrary: i j*)
 case 0
 then show ?case **using** *asm* **by** *simp*
next
 case (*Suc n*)
 hence *: $n = j - Suc \ i$ **by** *linarith*
 have $Suc \ i \leq j$ **using** *Suc(2,3)* **by** *linarith*
 thus ?case **using** *sets-F-mono*[*OF le-SucI*] *Suc(4)* *Suc(1)*[*OF **] **by** (*auto*)

intro: assms(2)[THEN order-trans])

qed

qed (*simp add: integrable*)

lemma (*in discrete-time-adapted-process-order*) *submartingale-nat*:

assumes *integrable*: $\bigwedge i. \text{integrable } M \ (X \ i)$

and $\bigwedge i. AE \ \xi \text{ in } M. X \ i \ \xi \leq \text{cond-exp } M \ (F \ i) \ (X \ (Suc \ i)) \ \xi$

shows *discrete-time-submartingale* $M \ F \ X$

using *subalg integrable assms(2)*

by (*intro submartingale-of-set-integral-le-Suc ord-le-eq-trans[OF set-integral-mono-AE-banach cond-exp-set-integral[symmetric]], simp*)

(*meson in-mono integrable-mult-indicator set-integrable-def subalgebra-def,*

meson integrable-cond-exp in-mono integrable-mult-indicator set-integrable-def

subalgebra-def,

auto simp add: subalgebra-def, metis (mono-tags, lifting) AE-I2 AE-mp)

lemma (*in discrete-time-adapted-process-order*) *submartingale-of-cond-exp-diff-Suc-nonneg*:

assumes *integrable*: $\bigwedge i. \text{integrable } M \ (X \ i)$

and $\bigwedge i. AE \ \xi \text{ in } M. 0 \leq \text{cond-exp } M \ (F \ i) \ (\lambda \xi. X \ (Suc \ i) \ \xi - X \ i \ \xi) \ \xi$

shows *discrete-time-submartingale* $M \ F \ X$

proof (*intro submartingale-nat integrable*)

fix i

show $AE \ \xi \text{ in } M. X \ i \ \xi \leq \text{cond-exp } M \ (F \ i) \ (X \ (Suc \ i)) \ \xi$ **using** *cond-exp-diff[OF integrable(1,1), of i Suc i i] cond-exp-F-meas[OF integrable adapted, of i] assms(2)[of i]* **by fastforce**

qed

5 Discrete Time Supermartingales

lemma (*in discrete-time-supermartingale*) *predictable-le-bot*:

assumes *predictable* X

shows $AE \ \xi \text{ in } M. X \ i \ \xi \leq X \ \perp \ \xi$

proof (*induction i*)

case 0

then show *?case* **by** (*simp add: bot-nat-def*)

next

case $(Suc \ i)$

thus *?case* **using** *predictable-discrete-time-process-measurable[OF assms, of Suc i]*

supermartingale-property[OF le-SucI, of i]

cond-exp-F-meas[OF integrable, of Suc i i] Suc **by fastforce**

qed

lemma (*in discrete-time-adapted-process-order*) *supermartingale-of-set-integral-ge-Suc*:

assumes *integrable*: $\bigwedge i. \text{integrable } M \ (X \ i)$

and $\bigwedge A \ i. A \in F \ i \implies \text{set-lebesgue-integral } M \ A \ (X \ (Suc \ i)) \leq \text{set-lebesgue-integral } M \ A \ (X \ i)$

shows *discrete-time-supermartingale* $M \ F \ X$

proof –

interpret *uminus-X: discrete-time-adapted-process-order* $M F - X$ **by** (*intro discrete-time-adapted-process-order.intro adapted-process-order.intro uminus*)
note $*$ = *set-integral-uminus*[*unfolded set-integrable-def*, *OF integrable-mult-indicator*[*OF - integrable*]]
have *discrete-time-supermartingale* $M F (-(- X))$ **using** *ord-eq-le-trans*[*OF * ord-le-eq-trans*[*OF le-imp-neg-le*[*OF assms*(2)] $*$ [*symmetric*]]] *subalg*
by (*intro discrete-time-supermartingale.intro submartingale.uminus discrete-time-submartingale.axioms uminus-X.submartingale-of-set-integral-le-Suc*) (*auto simp add: subalgebra-def integrable fun-Compl-def*, *blast*)
thus ?thesis **unfolding** *fun-Compl-def* **by** *simp*
qed

lemma (*in discrete-time-adapted-process-order*) *supermartingale-nat*:

assumes *integrable*: $\bigwedge i. \text{integrable } M (X i)$
and $\bigwedge i. AE \xi \text{ in } M. X i \xi \geq \text{cond-exp } M (F i) (X (Suc i)) \xi$
shows *discrete-time-supermartingale* $M F X$

proof –

interpret *uminus-X: discrete-time-adapted-process-order* $M F - X$ **by** (*intro discrete-time-adapted-process-order.intro adapted-process-order.intro uminus*)
have $AE \xi \text{ in } M. - X i \xi \leq \text{cond-exp } M (F i) (\lambda x. - X (Suc i) x) \xi$ **for** i **using** *assms*(2) *cond-exp-uminus*[*OF integrable*, *of i Suc i*] **by** *force*
hence *discrete-time-supermartingale* $M F (-(- X))$ **by** (*intro discrete-time-supermartingale.intro submartingale.uminus discrete-time-submartingale.axioms uminus-X.submartingale-nat*)
(*simp only: fun-Compl-def, intro integrable-minus integrable, auto simp add: fun-Compl-def*)
thus ?thesis **unfolding** *fun-Compl-def* **by** *simp*
qed

lemma (*in discrete-time-adapted-process-order*) *supermartingale-of-cond-exp-diff-Suc-nonneg*:

assumes *integrable*: $\bigwedge i. \text{integrable } M (X i)$
and $\bigwedge i. AE \xi \text{ in } M. 0 \leq \text{cond-exp } M (F i) (\lambda \xi. X i \xi - X (Suc i) \xi) \xi$
shows *discrete-time-supermartingale* $M F X$

proof (*intro supermartingale-nat integrable*)

fix i

show $AE \xi \text{ in } M. X i \xi \geq \text{cond-exp } M (F i) (X (Suc i)) \xi$ **using** *cond-exp-diff*[*OF integrable*(1,1), *of i i Suc i*] *cond-exp-F-meas*[*OF integrable adapted*, *of i*] *assms*(2)[*of i*] **by** *fastforce*

qed

end

References

- [1] D. Micciancio and S. Goldwasser. *Complexity of Lattice Problems: a Cryptographic Perspective*. Springer US, Boston, MA, 2002. OCLC: 852791069.