# On the Formalization of Martingales

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#### Abstract

In the scope of this project, we present a formalization of martingales in arbitrary Banach spaces using Isabelle/HOL [2].

The current formalization of conditional expectation in the Isabelle library is limited to real-valued functions. To overcome this limitation, we extend the construction of conditional expectation to general Banach spaces, employing an approach similar to the one described in [1]. We use measure theoretic arguments to construct the conditional expectation using suitable limits of simple functions.

Subsequently, we define stochastic processes and introduce the concepts of adapted, progressively measurable and predictable processes using suitable locale definitions<sup>1</sup>. We show the relation

 $adapted \supseteq progressive \supseteq predictable$ 

Furthermore, we show that progressive measurability and adaptedness are equivalent when the indexing set is discrete. We pay special attention to predictable processes in discrete-time, showing that  $(X_n)_{n\in\mathbb{N}}$  is predictable if and only if  $(X_{n+1})_{n\in\mathbb{N}}$  is adapted.

Moving forward, we rigorously define martingales, submartingales, and supermartingales, presenting their first consequences and corollaries<sup>2</sup>. Discrete-time martingales are given special attention in the formalization. In every step of our formalization, we make extensive use of the powerful locale system of Isabelle.

The formalization further contributes by generalizing concepts in Bochner integration by extending their application from the real numbers to arbitrary Banach spaces equipped with a second-countable topology. Induction schemes for integrable simple functions on Banach spaces are introduced, accommodating various scenarios with or without a real vector ordering<sup>3</sup>. Specifically, we formalize a powerful result called the "Averaging Theorem"[4] which allows us to show that densities are unique in Banach spaces.

In depth information on the formalization and the proofs of the individual theorems can be found in [3].

<sup>&</sup>lt;sup>1</sup>Martingale.Stochastic\_Process in [2]

<sup>&</sup>lt;sup>2</sup>Martingale.Martingale in [2]

 $<sup>^3</sup>$ Martingale.Bochner\_Integration\_Addendum in [2]

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8 Example: Coin Toss

```
theory Measure-Space-Supplement
imports HOL-Analysis.Measure-Space
begin
```

## 1 Supplementary Lemmas for Measure Spaces

## 1.1 $\sigma$ -Algebra Generated by a Family of Functions

```
definition family-vimage-algebra :: 'a set \Rightarrow ('a \Rightarrow 'b) set \Rightarrow 'b measure \Rightarrow 'a measure where
```

```
family-vimage-algebra \Omega S M \equiv sigma \Omega (\bigcup f \in S. \{f - `A \cap \Omega \mid A. A \in M\})
```

For singleton S, i.e.  $S = \{f\}$  for some f, the definition simplifies to that of vimage-algebra.

lemma family-vimage-algebra-singleton: family-vimage-algebra  $\Omega$  {f} M = vimage-algebra  $\Omega$  f M unfolding family-vimage-algebra-def vimage-algebra-def by simp

#### lemma

```
shows sets-family-vimage-algebra: sets (family-vimage-algebra \Omega S M) = sigma-sets \Omega (\bigcup f \in S. {f - 'A \cap \Omega \mid A. A \in M})
```

```
and space-family-vimage-algebra
[simp]: space (family-vimage-algebra \Omega 
 S M)=\Omega
```

by (auto simp add: family-vimage-algebra-def sets-measure-of-conv space-measure-of-conv)

 ${\bf lemma}\ measurable\mbox{-}family\mbox{-}vimage\mbox{-}algebra:$ 

```
assumes f \in S f \in \Omega \rightarrow space M
shows f \in family-vimage-algebra \Omega S M <math>\rightarrow_M M
using assms by (intro measurable I, auto simp add: sets-family-vimage-algebra)
```

 ${\bf lemma}\ measurable\mbox{-}family\mbox{-}vimage\mbox{-}algebra\mbox{-}singleton:$ 

```
assumes f \in \Omega \to space\ M
shows f \in family\text{-}vimage\text{-}algebra\ \Omega\ \{f\}\ M \to_M M
using assms measurable-family-vimage-algebra by blast
```

A collection of functions are measurable with respect to some  $\sigma$ -algebra N, if and only if the  $\sigma$ -algebra they generate is contained in N.

 $\mathbf{lemma}\ \textit{measurable-family-iff-sets}:$ 

```
shows (S \subseteq N \to_M M) \longleftrightarrow S \subseteq space N \to space M \land family-vimage-algebra (space N) <math>S M \subseteq N proof (standard, goal-cases)
```

hence subset:  $S \subseteq space \ N \rightarrow space \ M$  using measurable-space by fast

have  $\{f - `A \cap space \ N \mid A.\ A \in M\} \subseteq N \ \text{if} \ f \in S \ \text{for} \ f \ \text{using} \ measurable-iff-sets[unfolded family-vimage-algebra-singleton[symmetric], of f] 1 subset that by (fastforce simp add: sets-family-vimage-algebra)$ 

then show ?case unfolding sets-family-vimage-algebra using sets.sigma-algebra-axioms by (simp add: subset, intro sigma-algebra.sigma-sets-subset, blast+)

```
next
 case 2
 hence subset: S \subseteq space \ N \rightarrow space \ M by simp
 show ?case
 proof (standard, goal-cases)
   case (1 x)
   have family-vimage-algebra (space N) \{x\} M \subseteq N by (metis (no-types, lifting)
1 2 sets-family-vimage-algebra SUP-le-iff sigma-sets-le-sets-iff singletonD)
  thus ?case  using measurable - iff-sets[unfolded family-vimage-algebra-singleton[symmetric]]
subset[THEN subsetD, OF 1] by fast
 qed
qed
lemma family-vimage-algebra-diff:
 shows family-vimage-algebra \Omega S M = sigma \Omega (sets (family-vimage-algebra \Omega
(S-I) M) \cup family-vimage-algebra \Omega (S \cap I) M)
 using sets.space-closed space-measure-of-conv
 unfolding family-vimage-algebra-def sets-family-vimage-algebra
 by (intro sigma-eqI, blast, fastforce)
    (intro sigma-sets-eqI, blast, simp add: sets-measure-of-conv split: if-splits,
    meson\ Diff-subset Sup-subset-mono\ in-mono\ inf-sup-ord(1)\ sigma-sets-subseteq
subset-image-iff, fastforce+)
end
theory Elementary-Metric-Spaces-Supplement
 imports HOL-Analysis. Elementary-Metric-Spaces
begin
```

# 2 Supplementary Lemmas for Elementary Metric Spaces

## 2.1 Diameter Lemma

```
lemma diameter-comp-strict-mono: fixes s:: nat \Rightarrow 'a:: metric-space assumes strict-mono r bounded \{s\ i\ | i.\ r\ n \leq i\} shows diameter \{s\ (r\ i)\ |\ i.\ n \leq i\} \leq diameter\ \{s\ i\ |\ i.\ r\ n \leq i\} proof (rule\ diameter\text{-subset}) show \{s\ (r\ i)\ |\ i.\ n \leq i\} \subseteq \{s\ i\ |\ i.\ r\ n \leq i\} using assms(1) monotoneD strict-mono-mono by fastforce qed (intro\ assms(2))
lemma diameter\text{-bounded-bound'}: fixes S:: 'a:: metric\text{-space}\ set assumes S:\ bdd-above (case\text{-prod}\ dist\ '\ (S\times S))\ x \in S\ y \in S shows dist\ x\ y \leq diameter\ S proof - have (x,y) \in S\times S using S by auto then have dist\ x\ y \leq (SUP\ (x,y)\in S\times S.\ dist\ x\ y) by (rule\ cSUP\text{-upper2}]\ OF
```

```
assms(1)) simp
   with \langle x \in S \rangle show ?thesis by (auto simp: diameter-def)
qed
lemma bounded-imp-dist-bounded:
   assumes bounded (range s)
   shows bounded ((\lambda(i, j). dist (s i) (s j)) `(\{n..\} \times \{n..\}))
  using bounded-dist-comp[OF bounded-fst bounded-snd, OF bounded-Times(1,1)[OF
assms(1,1)] by (rule bounded-subset, force)
A sequence is Cauchy, if and only if it is bounded and it's diameter tends
to zero. The diameter is well-defined only if the sequence is bounded.
lemma cauchy-iff-diameter-tends-to-zero-and-bounded:
   fixes s :: nat \Rightarrow 'a :: metric\text{-}space
  shows Cauchy s \longleftrightarrow ((\lambda n. \ diameter \{ s \ i \mid i. \ i \geq n \}) \longrightarrow 0 \land bounded (range)
s))
proof -
   have \{s \ i \mid i. \ N \leq i\} \neq \{\} for N by blast
   hence diameter-SUP: diameter \{s \ i \ | i.\ N \leq i\} = (SUP\ (i,j) \in \{N..\} \times \{N..\}.
dist (s i) (s j)) for N unfolding diameter-def by (auto intro!: arg-cong[of - - Sup])
   show ?thesis
   proof (intro iffI)
      assume asm: Cauchy s
      have \exists N. \forall n \geq N. \text{ norm (diameter } \{s \ i \ | i. \ n \leq i\}) < e \ \text{if } e\text{-pos: } e > 0 \ \text{for } e
      proof -
           obtain N where dist-less: dist (s \ n) \ (s \ m) < (1/2) * e \ if \ n \ge N \ m \ge N
for n m using asm e-pos by (metis Cauchy-def mult-pos-pos zero-less-divide-iff
zero-less-numeral zero-less-one)
             fix r assume r \geq N
            hence dist (s \ n) \ (s \ m) < (1/2) * e \ \text{if} \ n \ge r \ m \ge r \ \text{for} \ n \ m \ \text{using} \ dist-less
that by simp
              hence (SUP\ (i, j) \in \{r..\} \times \{r..\}.\ dist\ (s\ i)\ (s\ j)) \le (1/2) * e by (intro\ intro\ intro
cSup-least) fastforce+
             also have \dots < e using e-pos by simp
           finally have diameter \{s \ i \ | i. \ r \leq i\} < e \ \text{using} \ diameter\text{-}SUP \ \text{by} \ presburger
         moreover have diameter \{s \mid i \mid i. r \leq i\} \geq 0 for r unfolding diameter-SUP
using bounded-imp-dist-bounded OF cauchy-imp-bounded, THEN bounded-imp-bdd-above,
OF \ asm] \ \mathbf{by} \ (intro \ cSup-upper2, \ auto)
          ultimately show ?thesis by auto
      qed
         thus (\lambda n. diameter \{s \ i \mid i. \ n < i\}) \longrightarrow 0 \land bounded (range s) using
cauchy-imp-bounded[OF asm] by (simp add: LIMSEQ-iff)
      assume asm: (\lambda n. \ diameter \{ s \ i \mid i. \ n \leq i \}) \longrightarrow 0 \land bounded (range s)
      have \exists N. \forall n \geq N. \forall m \geq N. dist(s n)(s m) < e \text{ if } e\text{-pos: } e > 0 \text{ for } e
      proof -
            obtain N where diam-less: diameter \{s \mid i \mid i. \ r \leq i\} < e \text{ if } r \geq N \text{ for } r
```

## 3 Supplementary Lemmas for Bochner Integration

## 3.1 Integrable Simple Functions

We restate some basic results concerning Bochner-integrable functions.

```
lemma integrable-implies-simple-function-sequence:
  fixes f :: 'a \Rightarrow 'b :: \{banach, second-countable-topology\}
  assumes integrable M f
  obtains s where \bigwedge i. simple-function M (s i)
     and \bigwedge i. emeasure M \{y \in space M. s i y \neq 0\} \neq \infty
     and \bigwedge x. \ x \in space \ M \Longrightarrow (\lambda i. \ s \ i \ x) \longrightarrow f \ x
     and \bigwedge x \ i. \ x \in space \ M \Longrightarrow norm \ (s \ i \ x) \le 2 * norm \ (f \ x)
proof-
  have f: f \in borel-measurable M (\int x, norm (f x) \partial M) < \infty using assms
unfolding integrable-iff-bounded by auto
  obtain s where s: \bigwedge i. simple-function M (s i) \bigwedge x. x \in space M \Longrightarrow (\lambda i. s
i \ x) \longrightarrow f \ x \ \land i \ x. \ x \in space \ M \Longrightarrow norm \ (s \ i \ x) \le 2 * norm \ (f \ x)  using
borel-measurable-implies-sequence-metric [OF \ f(1)] unfolding norm-conv-dist by
metis
  {
    have (\int x^+ \cdot x \cdot norm \ (s \ i \ x) \ \partial M) \le (\int x^+ \cdot x \cdot norm \ (f \ x) \cdot \partial M) using
s by (intro nn-integral-mono, auto)
  also have ... < \infty using f by (simp add: nn-integral-cmult enreal-mult-less-top
ennreal-mult)
    finally have sbi: Bochner-Integration.simple-bochner-integrable M (s i) using
s by (intro simple-bochner-integrable I-bounded) auto
     hence emeasure M \{y \in space M. \ s \ i \ y \neq 0\} \neq \infty by (auto intro: inte-
```

```
grable I-simple-bochner-integrable simple-bochner-integrable.cases)
   thus ?thesis using that s by blast
qed
Simple functions can be represented by sums of indicator functions.
lemma simple-function-indicator-representation:
    fixes f ::'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
   assumes simple-function M f x \in space M
   shows f x = (\sum y \in f \text{ 'space } M. \text{ indicator } (f - \{y\} \cap \text{space } M) \text{ } x *_R y)
    (is ? l = ? r)
proof -
   have ?r = (\sum y \in f \text{ 'space } M. (if y = f x \text{ then indicator } (f - `\{y\} \cap \text{space } M) \text{ } x *_R y \text{ else } \theta)) by (auto intro!:
   also have ... = indicator (f - `\{fx\} \cap space M) x *_R fx using assms by (auto
dest: simple-functionD)
   also have ... = f x using assms by (auto simp: indicator-def)
   finally show ?thesis by auto
qed
lemma simple-function-indicator-representation-AE:
   fixes f ::'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
   assumes simple-function Mf
   shows AE x in M. f x = (\sum y \in f \text{ 'space M. indicator } (f - `\{y\} \cap space M) x
     \mathbf{by} \ (\textit{metis} \ (\textit{mono-tags}, \ \textit{lifting}) \ \textit{AE-I2} \ \textit{simple-function-indicator-representation}
assms)
lemmas simple-function-scaleR[intro] = simple-function-compose2[\mathbf{where}\ h=(*_R)]
{\bf lemmas}\ integrable-simple-function = simple-bochner-integrable. intros [\it THEN\ has-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-integral-simple-bochner-inte
THEN integrable.intros
Induction rule for simple integrable functions.
lemma integrable-simple-function-induct[consumes 2, case-names cong indicator
add, induct set: simple-function]:
   fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
   assumes f: simple-function M f emeasure M \{y \in space M. f y \neq 0\} \neq \infty
   assumes cong: \bigwedge f q. simple-function M f \Longrightarrow emeasure M { y \in space\ M. f y \ne
\theta \} \neq \infty
                                        \implies simple-function M g \implies emeasure M \{y \in space M. g y \neq
\theta \} \neq \infty
    \Longrightarrow (\bigwedge x. \ x \in space \ M \Longrightarrow f \ x = g \ x) \Longrightarrow P \ f \Longrightarrow P \ g assumes indicator: \bigwedge A \ y. \ A \in sets \ M \Longrightarrow emeasure \ M \ A < \infty \Longrightarrow P \ (\lambda x.
indicator\ A\ x*_R\ y)
   assumes add: \bigwedge f g. simple-function M f \Longrightarrow emeasure M {y \in space\ M.\ f\ y \neq space\ M.
\theta \} \neq \infty \Longrightarrow
                                       simple-function\ M\ g \Longrightarrow emeasure\ M\ \{y \in space\ M.\ g\ y \neq 0\} \neq 0
\infty \Longrightarrow
```

```
(\bigwedge z. \ z \in space \ M \Longrightarrow norm \ (f \ z + g \ z) = norm \ (f \ z) + norm
(q z)) \Longrightarrow
                                            P f \Longrightarrow P g \Longrightarrow P (\lambda x. f x + g x)
   shows Pf
proof-
   let ?f = \lambda x. (\sum y \in f \text{ 'space } M. \text{ indicat-real } (f - `\{y\} \cap \text{ space } M) \ x *_R y)
   have f-ae-eq: fx = ?fx if x \in space M for x using simple-function-indicator-representation [OF]
   moreover have emeasure M {y \in space\ M. ?f\ y \neq 0} \neq \infty by (metis (no-types,
lifting) Collect-cong calculation f(2))
    moreover have P (\lambda x. \sum y \in S. indicat-real (f - `\{y\} \cap space M) \ x *_R y)
                                  simple-function M (\lambda x. \sum y \in S. indicat-real (f - `\{y\} \cap space M) x
*_R y)
                                   emeasure M \{ y \in space \ M. \ (\sum x \in S. \ indicat\text{-real} \ (f - `\{x\} \cap space \ f = `\{x
M) y *_R x) \neq 0 \} \neq \infty
                                   if S \subseteq f 'space M for S using simple-functionD(1)[OF\ assms(1),
 THEN rev-finite-subset, OF that that
   proof (induction rule: finite-induct)
        case empty
         {
            case 1
            then show ?case using indicator[of {}] by force
            case 2
            then show ?case by force
        next
            case 3
            then show ?case by force
        }
    next
        case (insert x F)
        have (f - `\{x\} \cap space M) \subseteq \{y \in space M. f y \neq 0\}  if x \neq 0 using that by
blast
        moreover have \{y \in space M. f y \neq 0\} = space M - (f - `\{0\} \cap space M)
           moreover have space M - (f - `\{0\} \cap space M) \in sets M  using sim-
ple-functionD(2)[OF f(1)] by blast
          ultimately have fin-0: emeasure M (f - (x) \cap space M) < \infty if x \neq 0
using that by (metis emeasure-mono f(2) infinity-enrical-def top.not-eq-extremum
top-unique)
      hence fin-1: emeasure M {y \in space M. indicat-real (f - `\{x\} \cap space M) y *_R 
x \neq 0} \neq \infty if x \neq 0 by (metis (mono-tags, lifting) emeasure-mono f(1) indica-
tor\text{-}simps(2) linorder\text{-}not\text{-}less mem\text{-}Collect\text{-}eq scaleR\text{-}eq\text{-}0\text{-}iff simple\text{-}functionD(2)
subsetI that)
      \mathbf{have} *: (\sum y \in insert \ x \ F. \ indicat\text{-}real \ (f \ -`\{y\} \cap space \ M) \ xa *_R y) = (\sum y \in F.
indicat-real (f - (y) \cap space M) xa *_R y) + indicat-real <math>(f - (x) \cap space M)
```

 $xa *_R x$  for xa by (metis (no-types, lifting) Diff-empty Diff-insert0 add.commute

insert.hyps(1) insert.hyps(2) sum.insert-remove)

```
have **: \{y \in space \ M. \ (\sum x \in insert \ x \ F. \ indicat\ real \ (f - `\{x\} \cap space \ M) \ y \}
*_R x) \neq 0\} \subseteq \{y \in space \ M. \ (\sum x \in F. \ indicat\ real \ (f - `\{x\} \cap space \ M) \ y *_R x)\}
\neq 0 \} \cup \{ y \in space M. indicat-real (f - `\{x\} \cap space M) \ y *_R x \neq 0 \} unfolding
* by fastforce
    {
     case 1
     hence x: x \in f 'space M and F: F \subseteq f 'space M by auto
     show ?case
     proof (cases \ x = \theta)
       case True
       then show ?thesis unfolding * using insert(3)[OF\ F] by simp
     \mathbf{next}
       case False
       have norm-argument: norm ((\sum y \in F. indicat\text{-real } (f - `\{y\} \cap space M) z))
*_R y) + indicat\text{-real } (f - `\{x\} \cap space M) \ z *_R x) = norm \ (\sum y \in F. indicat\text{-real})
(f - `\{y\} \cap space\ M)\ z *_R y) + norm\ (indicat-real\ (f - `\{x\} \cap space\ M)\ z *_R x)
if z: z \in space M for z
       proof (cases f z = x)
         case True
         have indicat-real (f - (y) \cap space M) z *_R y = 0 if y \in F for y using
True insert(2) z that 1 unfolding indicator-def by force
        hence (\sum y \in F. indicat\text{-}real (f - `\{y\} \cap space M) z *_R y) = 0 by (meson
sum.neutral)
         then show ?thesis by force
       next
         {\bf case}\ \mathit{False}
         then show ?thesis by force
       show ?thesis using False simple-functionD(2)[OF f(1)] insert(3,5)[OF F]
simple-function-scaleR fin-0 fin-1 by (subst *, subst add, subst simple-function-sum)
(blast\ intro:\ norm-argument\ indicator)+
     qed
   next
     case 2
     hence x: x \in f 'space M and F: F \subseteq f 'space M by auto
     show ?case
     proof (cases x = \theta)
       \mathbf{case} \ \mathit{True}
       then show ?thesis unfolding * using insert(4)[OF\ F] by simp
     next
       case False
     then show ?thesis unfolding * using insert(4)[OFF] simple-functionD(2)[OF]
f(1)] by fast
     qed
   \mathbf{next}
     hence x: x \in f 'space M and F: F \subseteq f 'space M by auto
     show ?case
     proof (cases x = \theta)
```

```
case True
             then show ?thesis unfolding * using insert(5)[OF\ F] by simp
          next
              case False
              have emeasure M \{y \in space M. (\sum x \in insert \ x \ F. \ indicat-real \ (f - `\{x\})\}
\cap space M) y *_R x) \neq 0} \leq emeasure M (\{y \in \text{space M. } (\sum x \in F. \text{ indicat-real } (f \in F) \})
-`\{x\} \cap space\ M)\ y *_R x) \neq 0\} \cup \{y \in space\ M.\ indicat\ real\ (f -`\{x\} \cap space\ M)\}
M) y *_{R} x \neq 0
             using ** simple-functionD(2)[OF\ insert(4)[OF\ F]]\ simple-functionD(2)[OF\ f]
f(1)] by (intro emeasure-mono, force+)
             also have ... \leq emeasure M \{y \in space M. (\sum x \in F. indicat\text{-real } (f - `\{x\})\}
\cap space M) y *_R x) \neq 0} + emeasure M \{y \in space M. indicat-real (f - `\{x\} \cap space M) \}
space M) y *_R x \neq \emptyset}
                 using simple-functionD(2)[OF\ insert(4)[OF\ F]]\ simple-functionD(2)[OF\ and below the context of the context
f(1)] by (intro emeasure-subadditive, force+)
               also have ... < \infty using insert(5)[OF F] fin-1[OF False] by (simp add:
less-top)
             finally show ?thesis by simp
          qed
      }
   qed
   moreover have simple-function M (\lambda x. \sum y \in f 'space M. indicat-real (f - `\{y\})
\cap space M) x *_R y) using calculation by blast
    moreover have P (\lambda x. \sum y \in f 'space M. indicat-real (f - `\{y\} \cap space M) x
*_R y) using calculation by blast
   ultimately show ?thesis by (intro cong[OF - - f(1,2)], blast, presburger+)
qed
Induction rule for non-negative simple integrable functions
lemma integrable-simple-function-induct-nn[consumes 3, case-names cong indica-
tor add, induct set: simple-function]:
    fixes f::'a \Rightarrow 'b:: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
   assumes f: simple-function M f emeasure M \{y \in space M. f y \neq 0\} \neq \infty \land x.
x \in space \ M \longrightarrow f \ x \ge 0
    assumes cong: \bigwedge f g. simple-function M f \Longrightarrow emeasure M {y \in space M. f y
\neq 0 \} \neq \infty \Longrightarrow (\bigwedge x. \ x \in space \ M \Longrightarrow f \ x \geq 0) \Longrightarrow simple-function \ M \ g \Longrightarrow
emeasure M \{ y \in space \ M. \ g \ y \neq 0 \} \neq \infty \Longrightarrow (\bigwedge x. \ x \in space \ M \Longrightarrow g \ x \geq 0 )
\Longrightarrow (\bigwedge x. \ x \in space \ M \Longrightarrow f \ x = g \ x) \Longrightarrow P \ f \Longrightarrow P \ g
  assumes indicator: \bigwedge A y. y \ge 0 \Longrightarrow A \in sets M \Longrightarrow emeasure M A < \infty \Longrightarrow
P(\lambda x. indicator A x *_R y)
    assumes add: \bigwedge f g. (\bigwedge x. \ x \in space \ M \Longrightarrow f \ x \geq 0) \Longrightarrow simple-function \ M \ f
\implies emeasure M \{ y \in space M. f y \neq 0 \} \neq \infty \implies
                                        (\bigwedge x.\ x\in \mathit{space}\ M\Longrightarrow g\ x\geq 0)\Longrightarrow \mathit{simple-function}\ M\ g\Longrightarrow
emeasure M \{ y \in space M. g y \neq 0 \} \neq \infty \Longrightarrow
                                       (\bigwedge z. \ z \in space \ M \Longrightarrow norm \ (f \ z + g \ z) = norm \ (f \ z) + norm
(g z)) \Longrightarrow
                                     P f \Longrightarrow P g \Longrightarrow P (\lambda x. f x + g x)
   shows Pf
```

```
let ?f = \lambda x. (\sum y \in f \text{ 'space } M. \text{ indicat-real } (f - `\{y\} \cap \text{space } M) \ x *_R y)
 have f-ae-eq: fx = ?fx if x \in space M for x using simple-function-indicator-representation [OF]
 moreover have emeasure M {y \in space\ M. ?f\ y \neq 0} \neq \infty by (metis (no-types,
lifting) Collect-cong calculation f(2))
 moreover have P (\lambda x. \sum y \in S. indicat-real (f - `\{y\} \cap space\ M)\ x *_R y) simple-function M (\lambda x. \sum y \in S. indicat-real (f - `\{y\} \cap space\ M)\ x
*_R y)
                 emeasure M \{y \in space M. (\sum x \in S. indicat\text{-real } (f - `\{x\} \cap space )\} \}
M) \ y *_R x) \neq \emptyset \} \neq \infty
               \bigwedge x. x \in space M \Longrightarrow 0 \le (\sum y \in S. indicat-real (f - `\{y\} \cap space M)
x *_R y
                 if S \subseteq f 'space M for S using simple-functionD(1)[OF \ assms(1),
THEN rev-finite-subset, OF that that
  proof (induction rule: finite-subset-induct')
    case empty
    {
      case 1
      then show ?case using indicator[of 0 \ \{\}] by force
      case 2
      then show ?case by force
    \mathbf{next}
      case 3
      then show ?case by force
    \mathbf{next}
      case 4
      then show ?case by force
    }
  \mathbf{next}
    case (insert x F)
    have (f - (x) \cap space M) \subseteq \{y \in space M. f y \neq 0\} if x \neq 0 using that by
    moreover have \{y \in space M. f y \neq 0\} = space M - (f - `\{0\} \cap space M)
     moreover have space M - (f - `\{\theta\} \cap space M) \in sets M using sim-
ple-functionD(2)[OF f(1)] by blast
     ultimately have fin-0: emeasure M (f - {}^{\cdot} \{x\} \cap space M) < \infty if x \neq 0
using that by (metis emeasure-mono f(2) infinity-enried-def top.not-eq-extremum
top-unique)
   hence fin-1: emeasure M {y \in space\ M. indicat-real (f - `\{x\} \cap space\ M)\ y *_R
x \neq 0 \neq \infty if x \neq 0 by (metis (mono-tags, lifting) emeasure-mono f(1) indica-
tor\text{-}simps(2)\ linorder\text{-}not\text{-}less\ mem\text{-}Collect\text{-}eq\ scaleR\text{-}eq\text{-}0\text{-}iff\ simple\text{-}function}D(2)
subsetI that)
    have nonneg-x: x \ge 0 using insert f by blast
have *: (\sum y \in insert \ x \ F. \ indicat-real \ (f \ -` \{y\} \cap space \ M) \ xa *_R y) = (\sum y \in F. \ indicat-real \ (f \ -` \{y\} \cap space \ M) \ xa *_R y) + indicat-real \ (f \ -` \{x\} \cap space \ M)
```

proof-

```
space M) xa *_R x for xa by (metis (no-types, lifting) add.commute insert.hyps(1)
insert.hyps(4) sum.insert)
   have **: \{y \in space \ M. \ (\sum x \in insert \ x \ F. \ indicat-real \ (f - `\{x\} \cap space \ M) \ y \}
*_R x) \neq 0\} \subseteq \{y \in space \ M. \ (\sum x \in F. \ indicat\ real \ (f - `\{x\} \cap space \ M) \ y *_R x)\}
\neq 0 \} \cup \{ y \in space \ M. \ indicat\ real \ (f - `\{x\} \cap space \ M) \ y *_R x \neq 0 \}  unfolding
* by fastforce
    {
     case 1
     show ?case
     proof (cases x = 0)
       case True
       then show ?thesis unfolding * using insert by simp
     next
       case False
       have norm-argument: norm ((\sum y \in F. indicat\text{-real } (f - `\{y\} \cap space M) z))
*_R y) + indicat\text{-real } (f - `\{x\} \cap space M) \ z *_R x) = norm \ (\sum y \in F. \ indicat\text{-real } )
(f - `\{y\} \cap space\ M)\ z *_R y) + norm\ (indicat-real\ (f - `\{x\} \cap space\ M)\ z *_R x)
if z: z \in space M for z
       proof (cases f z = x)
         case True
         have indicat-real (f - `\{y\} \cap space M) \ z *_R y = 0 \ \textbf{if} \ y \in F \ \textbf{for} \ y \ \textbf{using}
True insert z that 1 unfolding indicator-def by force
        hence (\sum y \in F. indicat\text{-}real (f - `\{y\} \cap space M) z *_R y) = 0 by (meson
sum.neutral)
         thus ?thesis by force
       qed (force)
      show ?thesis using False fin-0 fin-1 f norm-argument by (subst *, subst add,
presburger add: insert, intro insert, intro insert, fastforce simp add: indicator-def
intro!: insert(2) \ f(3), \ auto \ intro!: indicator \ insert \ nonneg-x)
     qed
   next
     case 2
     show ?case
     proof (cases x = \theta)
       \mathbf{case} \ \mathit{True}
       then show ?thesis unfolding * using insert by simp
     next
      then show ?thesis unfolding * using insert simple-functionD(2)[OF f(1)]
by fast
     qed
   next
     case \beta
     show ?case
     proof (cases x = \theta)
       case True
       then show ?thesis unfolding * using insert by simp
     next
       case False
```

```
have emeasure M \{y \in space M. (\sum x \in insert \ x \ F. \ indicat-real \ (f - `\{x\})\}
\cap space M) y *_R x) \neq 0 \leq emeasure M (\{y \in space M. (\sum x \in F. indicat-real (f \in Space M))\}
-`\{x\} \cap space\ M)\ y *_R x) \neq 0\} \cup \{y \in space\ M.\ indicat\ real\ (f -`\{x\} \cap space\ M)\}
M) y *_{R} x \neq 0
                 using ** simple-functionD(2)[OF\ insert(6)]\ simple-functionD(2)[OF\ f(1)]
insert.IH(2) by (intro emeasure-mono, blast, simp)
               also have ... \leq emeasure M \{y \in space M. (\sum x \in F. indicat\text{-real } (f - `\{x\})\}
\cap space M) y *_R x) \neq 0} + emeasure M \{y \in \text{space } M. \text{ indicat-real } (f - `\{x\} \cap X) \}
space M) y *_R x \neq 0}
                      using simple-functionD(2)[OF\ insert(6)]\ simple-functionD(2)[OF\ f(1)]
by (intro emeasure-subadditive, force+)
               also have ... < \infty using insert(7) fin-1[OF False] by (simp add: less-top)
                finally show ?thesis by simp
            qed
        next
            case (4 \xi)
        thus ?case using insert nonneq-x f(3) by (auto simp add: scaleR-nonneq-nonneq
intro: sum-nonneq)
        }
    qed
   moreover have simple-function M (\lambda x. \sum y \in f 'space M. indicat-real (f - `\{y\})
\cap space M) x *_R y) using calculation by blast
    moreover have P(\lambda x. \sum y \in f \text{ 'space } M. \text{ indicat-real } (f - \{y\} \cap \text{space } M) \text{ } x
*_R y) using calculation by blast
    moreover have \bigwedge x. x \in space M \Longrightarrow 0 \le f x using f(3) by simp
     ultimately show ?thesis by (intro cong[OF - - - f(1,2)], blast, blast, fast)
presburger +
qed
lemma finite-nn-integral-imp-ae-finite:
    fixes f :: 'a \Rightarrow ennreal
    assumes f \in borel-measurable M (\int x. f x \partial M) < \infty
    shows AE x in M. f x < \infty
proof (rule ccontr, goal-cases)
    case 1
    let ?A = space M \cap \{x. f x = \infty\}
    have *: emeasure M ?A > 0 using 1 assms(1) by (metis (mono-tags, lifting)
assms(2) eventually-mono infinity-ennreal-def nn-integral-noteq-infinite top.not-eq-extremum)
     (metis (mono-tags, lifting) indicator-inter-arith indicator-simps(2) mem-Collect-eq
mult-eq-\theta-iff)
  also have ... = \infty * emeasure M ?A  using assms(1) by (intro nn-integral-cmult-indicator,
simp)
    also have ... = \infty using * by fastforce
    finally have (\int x \cdot f \cdot x \cdot indicator ?A \cdot x \cdot \partial M) = \infty.
    moreover have (\int_{-\infty}^{+\infty} x \cdot f(x) \cdot f(x)
nn-integral-mono, simp add: indicator-def)
    ultimately show ?case using assms(2) by simp
qed
```

Convergence in L1-Norm implies existence of a subsequence which convergences almost everywhere. This lemma is easier to use than the existing one in HOL-Analysis.Bochner-Integration

```
lemma cauchy-L1-AE-cauchy-subseq:
  fixes s :: nat \Rightarrow 'a \Rightarrow 'b :: \{banach, second-countable-topology\}
  assumes [measurable]: \land n. integrable M (s n)
     and \bigwedge e. \ e > 0 \Longrightarrow \exists N. \ \forall i \geq N. \ \forall j \geq N. \ LINT \ x | M. \ norm \ (s \ i \ x - s \ j \ x) < e
  obtains r where strict-mono r AE x in M. Cauchy (\lambda i. s (r i) x)
  have \exists r. \forall n. (\forall i \geq r \ n. \ \forall j \geq r \ n. \ LINT \ x | M. \ norm \ (s \ i \ x - s \ j \ x) < (1 \ / \ 2) 
n) \wedge (r (Suc \ n) > r \ n)
  proof (intro dependent-nat-choice, goal-cases)
   case 1
   then show ?case using assms(2) by presburger
  next
   case (2 x n)
   obtain N where *: LINT x|M. norm (s i x - s j x) < (1 / 2) \cap Suc n if i \ge
N j \ge N for i j using assms(2)[of (1 / 2) \cap Suc n] by auto
     fix i j assume i \ge max \ N \ (Suc \ x) \ j \ge max \ N \ (Suc \ x)
     hence integral^L M (\lambda x. norm (s i x - s j x)) < (1 / 2) ^Suc n using * by
   then show ?case by fastforce
  qed
  then obtain r where strict-mono: strict-mono r and \forall i \geq r \ n. \ \forall j \geq r \ n. \ LINT
x|M. norm (s \ i \ x - s \ j \ x) < (1 \ / \ 2) \ \hat{} \ n for n using strict-mono-Suc-iff by blast
  hence r-is: LINT x|M. norm (s(r(Suc n)) x - s(r n) x) < (1/2) ^n for n
by (simp add: strict-mono-leD)
  define g where g = (\lambda n \ x. \ (\sum i \le n. \ ennreal \ (norm \ (s \ (r \ (Suc \ i)) \ x - s \ (r \ i))))
  define g' where g' = (\lambda x. \sum i. ennreal (norm (s (r (Suc i)) x - s (r i) x)))
  have integrable-g: (\int + x. g n x \partial M) < 2 for n
  proof -
    have (\int x. g \, n \, x \, \partial M) = (\int x. (\sum i \leq n. ennreal (norm (s (r (Suc i)) x - i)))
s\ (r\ i)\ x)))\ \partial M) using g-def by simp
    also have ... = (\sum_{i \le n} i \le n). (\int_{-\infty}^{\infty} x)^{-1} = (\sum_{i \le n} i \le n). (f(x) = \sum_{i \le n} i \le n).
\partial M)) by (intro nn-integral-sum, simp)
     also have ... = (\sum i \le n. LINT x|M. norm (s (r (Suc i)) x - s (r i) x))
unfolding dist-norm using assms(1) by (subst nn-integral-eq-integral[OF inte-
grable-norm], auto)
   also have ... < ennreal (\sum i \le n. (1/2) \hat{i}) by (intro ennreal-lessI[OF sum-pos
sum-strict-mono[OF finite-atMost - r-is]], auto)
   also have ... \leq ennreal\ 2 unfolding sum\text{-}gp0[of\ 1\ /\ 2\ n] by (intro ennreal-leI,
   finally show (\int x \cdot g \, n \, x \, \partial M) < 2 by simp
  qed
```

```
have integrable-g': (\int x \cdot g' x \cdot \partial M) \leq 2
  proof -
     have incseq (\lambda n. \ g \ n \ x) for x by (intro incseq-SucI, auto simp add: g-def
ennreal-leI)
     hence convergent (\lambda n. \ g \ n \ x) for x unfolding convergent-def using LIM-
SEQ-incseq-SUP by fast
      hence (\lambda n. \ g \ n \ x) \longrightarrow g' \ x for x unfolding g-def g'-def by (intro
summable-iff-convergent'[THEN iffD2, THEN summable-LIMSEQ'], blast)
  hence (\int x, g' x \partial M) = (\int x, liminf(\lambda n, g n x) \partial M) by (metis lim-imp-Liminf)
trivial-limit-sequentially)
   also have ... \leq liminf(\lambda n. \int + x. g n x \partial M) by (intro nn-integral-liminf, simp
add: g\text{-}def)
   also have ... \leq liminf(\lambda n. 2) using integrable-g by (intro\ Liminf-mono) (simp
add: order-le-less)
   also have ... = 2 using sequentially-bot tendsto-iff-Liminf-eq-Limsup by blast
   finally show ?thesis.
  qed
 hence AE x in M. g' x < \infty by (intro finite-nn-integral-imp-ae-finite) (auto simp
add: order-le-less-trans g'-def)
  moreover have summable (\lambda i. \ norm \ (s \ (r \ (Suc \ i)) \ x - s \ (r \ i) \ x)) if g' \ x \neq \infty
for x using that unfolding g'-def by (intro summable-suminf-not-top) fastforce+
  ultimately have ae-summable: AE x in M. summable (\lambda i.\ s\ (r\ (Suc\ i))\ x-s
(r i) x) using summable-norm-cancel unfolding dist-norm by force
   fix x assume summable (\lambda i. s (r (Suc i)) x - s (r i) x)
   hence (\lambda n. \sum i < n. \ s \ (r \ (Suc \ i)) \ x - s \ (r \ i) \ x) \longrightarrow (\sum i. \ s \ (r \ (Suc \ i)) \ x - s \ (r \ i) \ x)
s\ (r\ i)\ x)\ \mathbf{using}\ summable\text{-}LIMSEQ\ \mathbf{by}\ blast
   \mathbf{moreover\ have}\ (\lambda n.\ (\sum i < n.\ s\ (r\ (Suc\ i))\ x\ -\ s\ (r\ i)\ x)) = (\lambda n.\ s\ (r\ n)\ x\ -\ n)
s\ (r\ \theta)\ x) using sum\text{-}lessThan\text{-}telescope} by fastforce
   ultimately have (\lambda n. \ s \ (r \ n) \ x - s \ (r \ 0) \ x) \longrightarrow (\sum i. \ s \ (r \ (Suc \ i)) \ x - s)
(r \ i) \ x) by argo
   hence (\lambda n. \ s \ (r \ n) \ x - s \ (r \ 0) \ x + s \ (r \ 0) \ x) \longrightarrow (\sum i. \ s \ (r \ (Suc \ i)) \ x - s
(r \ i) \ x) + s \ (r \ \theta) \ x \ by \ (intro \ isCont-tends to-compose [of - \lambda z. \ z + s \ (r \ \theta) \ x], \ auto)
   hence Cauchy (\lambda n.\ s\ (r\ n)\ x) by (simp\ add:\ LIMSEQ-imp-Cauchy)
  hence AE x in M. Cauchy (\lambda i.\ s\ (r\ i)\ x) using ae-summable by fast
  thus ?thesis by (rule\ that[OF\ strict-mono(1)])
qed
3.2
        Totally Ordered Banach Spaces
lemma integrable-max[simp, intro]:
  \mathbf{fixes}\ f :: \ 'a \Rightarrow \ 'b :: \{second\text{-}countable\text{-}topology,\ banach,\ linorder\text{-}topology}\}
  assumes fg[measurable]: integrable M f integrable M g
  shows integrable M (\lambda x. max (f x) (g x))
proof (rule Bochner-Integration.integrable-bound)
```

```
fix x y :: 'b
    have norm (max \ x \ y) \le max \ (norm \ x) \ (norm \ y) by linarith
    also have ... \leq norm \ x + norm \ y \ by \ simp
    finally have norm (max \ x \ y) \le norm (norm \ x + norm \ y) by simp
 thus AE \ x \ in \ M. \ norm \ (max \ (f \ x) \ (g \ x)) \leq norm \ (norm \ (f \ x) + norm \ (g \ x)) by
qed (auto intro: Bochner-Integration.integrable-add[OF integrable-norm[OF fq(1)]
integrable-norm[OF\ fg(2)]])
lemma integrable-min[simp, intro]:
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology\}
  assumes [measurable]: integrable M f integrable M g
 shows integrable M (\lambda x. min (f x) (q x))
proof -
 have norm (min (f x) (g x)) \le norm (f x) \lor norm (min (f x) (g x)) \le norm (g x)
x) for x by linarith
 thus ?thesis by (intro integrable-bound OF integrable-max OF integrable-norm (1,1),
OF \ assms]], \ auto)
qed
Restatement of integral-nonneg-AE for functions taking values in a Banach
space.
lemma integral-nonneg-AE-banach:
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes [measurable]: f \in borel-measurable M and nonneg: AE \times in M. 0 \le f \times in M
  shows 0 \leq integral^L M f
proof cases
  assume integrable: integrable M f
  hence max: (\lambda x. \ max \ \theta \ (f \ x)) \in borel-measurable \ M \ \land x. \ \theta \leq max \ \theta \ (f \ x)
integrable M (\lambda x. max \theta (f x)) by auto
  hence 0 \leq integral^L M (\lambda x. max 0 (f x))
  proof -
  obtain s where *: \bigwedge i. simple-function M (s i)
                    \bigwedge i. emeasure M \{ y \in space M. \ s \ i \ y \neq 0 \} \neq \infty
                    \bigwedge x. \ x \in space \ M \Longrightarrow (\lambda i. \ s \ i \ x) \longrightarrow f \ x
                      \bigwedge x \ i. \ x \in space \ M \Longrightarrow norm \ (s \ i \ x) \le 2 * norm \ (f \ x) \ using
integrable-implies-simple-function-sequence[OF integrable] by blast
    have simple: \bigwedge i. simple-function M (\lambda x. max \theta (s i x)) using * by fast
    have \Lambda i. \{y \in space M. max \theta (s i y) \neq \theta\} \subseteq \{y \in space M. s i y \neq \theta\}
unfolding max-def by force
   \mathbf{moreover} \ \mathbf{have} \ \big \langle i. \ \{y \in \mathit{space} \ M. \ \mathit{s} \ i \ y \neq \emptyset \} \in \mathit{sets} \ M \ \mathbf{using} * \mathbf{by} \ \mathit{measurable}
     ultimately have \bigwedge i. emeasure M \{y \in space M. max 0 (s i y) \neq 0\} \leq
emeasure M \{ y \in space \ M. \ s \ i \ y \neq 0 \} by (rule emeasure-mono)
    hence **:\bigwedge i. emeasure M \{ y \in space M. max 0 (s i y) \neq 0 \} \neq \infty  using *(2)
by (auto intro: order.strict-trans1 simp add: top.not-eq-extremum)
    have \bigwedge x. \ x \in space \ M \Longrightarrow (\lambda i. \ max \ \theta \ (s \ i \ x)) \longrightarrow max \ \theta \ (f \ x) \ using *(3)
```

```
tendsto-max by blast
    moreover have \bigwedge x \ i. \ x \in space \ M \Longrightarrow norm \ (max \ 0 \ (s \ i \ x)) \leq norm \ (2 *_R)
f(x) using *(4) unfolding max-def by auto
  ultimately have tendsto: (\lambda i. integral^L M (\lambda x. max \theta (s i x))) \longrightarrow integral^L
M (\lambda x. max \theta (f x))
        using borel-measurable-simple-function simple integrable by (intro inte-
qral-dominated-convergence[OF\ max(1)\ -\ integrable-norm[OF\ integrable-scaleR-right],
of - 2f, blast+)
   {
      \mathbf{fix}\ h\ ::\ 'a\Rightarrow\ 'b::\{second\text{-}countable\text{-}topology,\ banach,\ linorder\text{-}topology,\ or\text{-}
dered-real-vector}
     assume simple-function M h emeasure M \{y \in space M. h y \neq 0\} \neq \infty \land x.
x \in space \ M \longrightarrow h \ x \ge 0
     \mathbf{hence} \, *: \, integral^L \,\, M \,\, h \, \geq \, \theta
     \mathbf{proof} (induct rule: integrable-simple-function-induct-nn)
       case (conq f q)
       then show ?case using Bochner-Integration.integral-cong by force
     next
       case (indicator A y)
       hence A \neq \{\} \implies y \geq 0 using sets.sets-into-space by fastforce
          then show ?case using indicator by (cases A = \{\}, auto simp add:
scaleR-nonneg-nonneg)
     next
       case (add f g)
       then show ?case by (simp add: integrable-simple-function)
   }
   thus ?thesis using LIMSEQ-le-const[OF tendsto, of 0] ** simple by fastforce
 also have ... = integral^L M f using nonneg by (auto intro: integral-cong-AE)
 finally show ?thesis.
qed (simp add: not-integrable-integral-eq)
lemma integral-mono-AE-banach:
  fixes f g :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes integrable M f integrable M g AE x in M. f x \leq g x
 shows integral^L M f < integral^L M g
 using integral-nonneg-AE-banach[of \ \lambda x. \ g \ x - f \ x] assms Bochner-Integration.integral-diff[OF]
assms(1,2)] by force
lemma integral-mono-banach:
  fixes f g :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes integrable M f integrable M g \bigwedge x. x \in space M \Longrightarrow f x \leq g x
 shows integral^L M f \leq integral^L M g
  using integral-mono-AE-banach assms by blast
```

## 3.3 Integrability and Measurability of the Diameter

```
context
  fixes s:: nat \Rightarrow 'a \Rightarrow 'b:: \{second\text{-}countable\text{-}topology, banach} \} and M
  assumes bounded: \bigwedge x. \ x \in space \ M \Longrightarrow bounded \ (range \ (\lambda i. \ s \ i \ x))
begin
{\bf lemma}\ borel-measurable-diameter:
  assumes [measurable]: \bigwedge i. (s i) \in borel-measurable M
 shows (\lambda x. \ diameter \{ s \ i \ x \ | i. \ n \leq i \}) \in borel-measurable M
proof -
  have \{s \ i \ x \mid i.\ N < i\} \neq \{\} for x \ N by blast
 hence diameter-SUP: diameter \{s \mid i \mid i \mid N \leq i\} = (SUP(i, j) \in \{N..\} \times \{N..\}).
dist\ (s\ i\ x)\ (s\ j\ x)) for x\ N unfolding diameter-def by (auto\ intro!:\ arg-cong[of\ -
- Sup
 have case-prod dist '(\{s \ i \ x \ | i. \ n \leq i\} \times \{s \ i \ x \ | i. \ n \leq i\}) = ((\lambda(i, j). dist (s \ i
(s \ j \ x) \ (\{n..\} \times \{n..\})) for x \ by \ fast
  hence *: (\lambda x. \ diameter \{s \ i \ x \ | i. \ n \le i\}) = (\lambda x. \ Sup \ ((\lambda(i, j). \ dist \ (s \ i \ x) \ (s \ j)) \}
(n...) \cdot (n...) \times (n...) using diameter-SUP by (simp add: case-prod-beta')
 have bounded ((\lambda(i, j). dist (s i x) (s j x)) `(\{n..\} \times \{n..\})) if x \in space M for
x by (rule bounded-imp-dist-bounded[OF bounded, OF that])
 hence bdd: bdd-above ((\lambda(i,j), dist (s i x) (s j x)) (\{n..\} \times \{n..\})) if x \in space
M for x using that bounded-imp-bdd-above by presburger
 have fst p \in borel-measurable M snd p \in borel-measurable M if p \in s '\{n..\} \times
s '\{n..\} for p using that by fastforce+
 hence (\lambda x. fst \ p \ x - snd \ p \ x) \in borel-measurable M \ \textbf{if} \ p \in s \ (n.) \times s \ (n.)
for p using that borel-measurable-diff by simp
  hence (\lambda x. \ case \ p \ of \ (f, \ g) \Rightarrow \ dist \ (f \ x) \ (g \ x)) \in borel-measurable \ M \ \textbf{if} \ p \in s \ `
\{n..\} \times s '\{n..\} for p unfolding dist-norm using that by measurable
  moreover have countable (s '\{n..\} × s '\{n..\}) by (intro countable-SIGMA
countable-image, auto)
  ultimately show ?thesis unfolding * by (auto intro!: borel-measurable-cSUP
bdd)
qed
lemma integrable-bound-diameter:
  fixes f :: 'a \Rightarrow real
  assumes integrable M f
      and [measurable]: \land i. (s i) \in borel-measurable M
      and \bigwedge x \ i. \ x \in space \ M \Longrightarrow norm \ (s \ i \ x) \le f \ x
    shows integrable M (\lambda x. diameter \{s \ i \ x \ | i. \ n \leq i\})
proof -
  have \{s \ i \ x \mid i.\ N \leq i\} \neq \{\} for x \ N by blast
 hence diameter-SUP: diameter \{s \ i \ x \ | i.\ N \leq i\} = (SUP\ (i,j) \in \{N..\} \times \{N..\}.
dist (s i x) (s j x)) for x N unfolding diameter-def by (auto intro!: arg-cong[of -
- Sup])
  {
    fix x assume x: x \in space M
```

```
let ?S = (\lambda(i, j). \ dist \ (s \ i \ x) \ (s \ j \ x)) \ `(\{n..\} \times \{n..\})
           have case-prod dist '(\{s \ i \ x \ | i. \ n \leq i\}) \times \{s \ i \ x \ | i. \ n \leq i\}) = (\lambda(i, j). \ dist \ (s \ i \ j) = (\lambda(i, j)) + (\lambda(i, j))
(s \ j \ x)) '(\{n..\} \times \{n..\}) by fast
           hence *: diameter \{s \ i \ x \ | i. \ n \leq i\} = Sup \ ?S \ using \ diameter-SUP \ by \ (simp)
add: case-prod-beta')
           have bounded ?S by (rule bounded-imp-dist-bounded[OF bounded[OF x]])
       hence Sup-S-nonneg: 0 \le Sup ?S by (auto intro!: cSup-upper2 x bounded-imp-bdd-above)
              have dist (s \ i \ x) \ (s \ j \ x) \le 2 * f \ x for i \ j \ by \ (intro \ dist-triangle 2 \ | THEN
order-trans, of - 0) (metis norm-conv-dist \ assms(3) \ x \ add-mono \ mult-2)
           hence \forall c \in ?S. \ c \leq 2 * f x \text{ by } force
           hence Sup ?S \le 2 * fx by (intro\ cSup\ least,\ auto)
           hence norm (Sup ?S) \le 2 * norm (f x) using Sup-S-nonneg by auto
           also have ... = norm (2 *_R f x) by simp
           finally have norm (diameter \{s \ i \ x \ | i. \ n \leq i\}) \leq norm \ (2 *_R f x) unfolding
    hence AE \ x \ in \ M. \ norm \ (diameter \ \{s \ i \ x \ | i. \ n \leq i\}) \leq norm \ (2 *_R f \ x) by blast
    thus integrable M (\lambda x. diameter {s \ i \ x \mid i.\ n \leq i}) using borel-measurable-diameter
by (intro\ Bochner-Integration.integrable-bound[OF\ assms(1)]THEN\ integrable-scaleR-right[of\ Assms(1)]THEN\
2]]], measurable)
qed
end
                         Auxiliary Lemmas for Set Integrals
3.4
\mathbf{lemma} set	ent-integral	ent-scaleR	ent-left:
      assumes A \in sets \ M \ c \neq 0 \Longrightarrow integrable \ M \ f
      shows LINT t:A|M. f t *_R c = (LINT t:A|M. f t) *_R c
      unfolding set-lebesgue-integral-def
      using integrable-mult-indicator[OF assms]
      by (subst integral-scaleR-left[symmetric], auto)
lemma nn-set-integral-eq-set-integral:
      assumes [measurable]: integrable M f
                and AE x \in A in M. 0 \le f x A \in sets M
           shows (\int x \in A \cdot f x \partial M) = (\int x \in A \cdot f x \partial M)
proof-
      have (\int x \cdot indicator A \cdot x *_R f \cdot x \cdot \partial M) = (\int x \in A \cdot f \cdot x \cdot \partial M)
    unfolding set-lebesque-integral-def using assms(2) by (intro nn-integral-eq-integral of
- \lambda x. indicat-real A x *_R f x], blast intro: assms integrable-mult-indicator, fastforce)
    moreover have (\int_{-\infty}^{\infty} x \cdot indicator A \times_R f \times \partial M) = (\int_{-\infty}^{\infty} x \in A \cdot f \times \partial M) by (metis
ennreal-0 indicator-simps(1) indicator-simps(2) mult.commute mult-1 mult-zero-left
real-scaleR-def)
      ultimately show ?thesis by argo
qed
```

**lemma** set-integral-restrict-space:

```
fixes f :: 'a \Rightarrow 'b :: \{banach, second-countable-topology\}
 assumes \Omega \cap space M \in sets M
 shows set-lebesgue-integral (restrict-space M \Omega) A f = set-lebesgue-integral M A
(\lambda x. indicator \Omega x *_R f x)
 unfolding set-lebesque-integral-def
 by (subst integral-restrict-space, auto intro!: integrable-mult-indicator assms simp:
mult.commute)
lemma set-integral-const:
  fixes c :: 'b :: \{banach, second\text{-}countable\text{-}topology\}
 assumes A \in sets \ M \ emeasure \ M \ A \neq \infty
 shows set-lebesgue-integral M A (\lambda-. c) = measure M A *_R c
 unfolding set-lebesgue-integral-def
 using assms by (metis has-bochner-integral-indicator has-bochner-integral-integral-eq
infinity-ennreal-def less-top)
lemma set-integral-mono-banach:
  fixes fg: 'a \Rightarrow 'b: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes set-integrable M A f set-integrable M A g
    \bigwedge x. \ x \in A \Longrightarrow f \ x \leq g \ x
 shows (LINT x:A|M. f x) \le (LINT x:A|M. g x)
  using assms unfolding set-integrable-def set-lebesgue-integral-def
 by (auto intro: integral-mono-banach split: split-indicator)
lemma set-integral-mono-AE-banach:
  fixes f g :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes set-integrable M A f set-integrable M A g AE x \in A in M. f x \leq g x
  shows set-lebesgue-integral M A f \leq set-lebesgue-integral M A g using assms
unfolding set-lebesgue-integral-def by (auto simp add: set-integrable-def intro!:
integral-mono-AE-banach[of\ M\ \lambda x.\ indicator\ A\ x*_R fx\ \lambda x.\ indicator\ A\ x*_R g\ x],
simp add: indicator-def)
```

### 3.5 Averaging Theorem

We aim to lift results from the real case to arbitrary Banach spaces. Our fundamental tool in this regard will be the averaging theorem. The proof of this theorem is due to Serge Lang (Real and Functional Analysis). The theorem allows us to make statements about a functions value almost everywhere, depending on the value its integral takes on various sets of the measure space.

Before we introduce and prove the averaging theorem, we will first show the following lemma which is crucial for our proof. While not stated exactly in this manner, our proof makes use of the characterization of second-countable topological spaces given in the book General Topology by Ryszard Engelking (Theorem 4.1.15).

```
lemma balls-countable-basis:
  obtains D :: 'a :: \{metric\text{-}space, second\text{-}countable\text{-}topology}\} set
  where topological-basis (case-prod ball '(D \times (\mathbb{Q} \cap \{0 < ..\})))
   and countable D
   and D \neq \{\}
proof -
  obtain D: 'a set where dense-subset: countable D D \neq \{\} [open U; U \neq \{\}]
\implies \exists y \in D. \ y \in U \text{ for } U \text{ using } countable\text{-}dense\text{-}exists \text{ by } blast
  have topological-basis (case-prod ball '(D \times (\mathbb{Q} \cap \{0 < ... \})))
 proof (intro topological-basis-iff[THEN iffD2], fast, clarify)
   fix U and x :: 'a assume asm: open U x \in U
   obtain e where e: e > 0 ball x \in U using asm openE by blast
  obtain y where y: y \in D y \in ball x (e / 3) using dense-subset(3)[OF open-ball,
of x \in /3 centre-in-ball [THEN iffD2, OF divide-pos-pos[OF e(1), of 3]] by force
  obtain r where r: r \in \mathbb{Q} \cap \{e/3 < ... < e/2\} unfolding Rats-def using of-rat-dense OF
divide-strict-left-mono[OF - e(1)], of 2 3 ] by auto
   have *: x \in ball \ y \ r \ using \ r \ y \ by \ (simp \ add: \ dist-commute)
   hence ball y r \subseteq U using r by (intro order-trans[OF - e(2)], simp, metric)
    moreover have ball y \in (case-prod\ ball\ (D \times (\mathbb{Q} \cap \{\theta < ..\}))) using y(1)
   ultimately show \exists B' \in (case\text{-prod ball } (D \times (\mathbb{Q} \cap \{0 < ..\}))). \ x \in B' \wedge B' \subseteq
U using * by meson
 qed
  thus ?thesis using that dense-subset by blast
qed
context sigma-finite-measure
begin
To show statements concerning \sigma-finite measure spaces, one usually shows
the statement for finite measure spaces and uses a limiting argument to show
it for the \sigma-finite case. The following induction scheme formalizes this.
lemma sigma-finite-measure-induct[case-names finite-measure, consumes 0]:
 assumes \bigwedge(N :: 'a \ measure) \ \Omega. finite-measure N
                            \implies N = restrict\text{-}space \ M \ \Omega
                            \Longrightarrow \Omega \in sets M
                            \implies emeasure\ N\ \Omega \neq \infty
                            \implies emeasure \ N \ \Omega \neq 0
                            \implies almost-everywhere N Q
     and [measurable]: Measurable.pred M Q
 shows almost-everywhere M Q
  have *: almost-everywhere N Q if finite-measure N N = restrict-space M \Omega \Omega
\in sets M emeasure N \Omega \neq \infty for N \Omega using that by (cases emeasure N \Omega = 0,
auto intro: emeasure-0-AE \ assms(1))
 obtain A:: nat \Rightarrow 'a \text{ set where } A: range A \subseteq sets M (\bigcup i. A i) = space M \text{ and}
emeasure-finite: emeasure M (A i) \neq \infty for i using sigma-finite by metis
```

```
note A(1)[measurable]
 have space-restr: space (restrict-space M(A i)) = A i for i unfolding space-restrict-space
by simp
  {
   \mathbf{fix} i
   have *: \{x \in A \ i \cap space \ M. \ Q \ x\} = \{x \in space \ M. \ Q \ x\} \cap (A \ i) by fast
  have Measurable.pred (restrict-space M(A i)) Q using A by (intro measurable I,
auto simp add: space-restr intro!: sets-restrict-space-iff[THEN iffD2], measurable,
auto)
  }
 note this[measurable]
  {
   \mathbf{fix} i
   have finite-measure (restrict-space M (A i)) using emeasure-finite by (intro
finite-measureI, subst space-restr, subst emeasure-restrict-space, auto)
   hence emeasure (restrict-space M (A i)) \{x \in A : \neg Q : x\} = 0 using emea-
sure-finite by (intro AE-iff-measurable THEN iffD1, OF - - *], measurable, subst
space-restr[symmetric], intro sets.top, auto simp add: emeasure-restrict-space)
  hence emeasure M \{x \in A \ i. \ \neg Q \ x\} = 0 by (subst emeasure-restrict-space[symmetric],
auto)
  hence emeasure M (\bigcup i. \{x \in A \ i. \ \neg \ Q \ x\}) = \theta by (intro emeasure-UN-eq-\theta,
  moreover have (\bigcup i. \{x \in A \ i. \ \neg \ Q \ x\}) = \{x \in space \ M. \ \neg \ Q \ x\} \text{ using } A \text{ by }
auto
  ultimately show ?thesis by (intro AE-iff-measurable[THEN iffD2], auto)
qed
The Averaging Theorem allows us to make statements concerning how a
function behaves almost everywhere, depending on its behaviour on average.
lemma averaging-theorem:
  fixes f::- \Rightarrow 'b::\{second\text{-}countable\text{-}topology, banach\}
 assumes [measurable]: integrable M f
     and closed: closed S
      and \bigwedge A. A \in sets \ M \Longrightarrow measure \ M \ A > 0 \Longrightarrow (1 \ / measure \ M \ A) *_R
set-lebesque-integral M A f \in S
   shows AE \ x \ in \ M. \ f \ x \in S
proof (induct rule: sigma-finite-measure-induct)
  case (finite-measure N \Omega)
 interpret finite-measure N by (rule finite-measure)
 have integrable measurable: integrable N f using assms finite-measure by (auto
simp: integrable-restrict-space integrable-mult-indicator)
  have average: (1 / Sigma-Algebra.measure N A) *<sub>R</sub> set-lebesgue-integral N A f
\in S \text{ if } A \in sets \ N \ measure \ N \ A > 0 \ \text{for } A
  have *: A \in sets M using that by (simp add: sets-restrict-space-iff finite-measure)
   have A = A \cap \Omega by (metis finite-measure(2,3) inf.orderE sets.sets-into-space
```

```
space-restrict-space that(1)
    hence set-lebesgue-integral N A f = set-lebesgue-integral M A f unfolding
finite-measure by (subst set-integral-restrict-space, auto simp add: finite-measure
set-lebesgue-integral-def indicator-inter-arith[symmetric])
    moreover have measure N A = measure M A using that by (auto intro!:
measure-restrict-space simp add: finite-measure sets-restrict-space-iff)
   ultimately show ?thesis using that * assms(3) by presburger
 qed
 obtain D: 'b set where balls-basis: topological-basis (case-prod ball '(D \times (\mathbb{Q}
\cap \{0 < ... \})) and countable-D: countable D using balls-countable-basis by blast
  have countable-balls: countable (case-prod ball ' (D \times (\mathbb{Q} \cap \{\theta < ...\}))) using
countable-rat countable-D by blast
 obtain B where B-balls: B \subseteq case\text{-prod ball} \ (D \times (\mathbb{Q} \cap \{0 < ..\})) \mid B = -S
using topological-basis[THEN iffD1, OF balls-basis] open-Compl[OF assms(2)] by
meson
 hence countable-B: countable B using countable-balls countable-subset by fast
 define b where b = from\text{-}nat\text{-}into\ (B \cup \{\{\}\}\})
 have B \cup \{\{\}\} \neq \{\} by simp
 have range-b: range b = B \cup \{\{\}\} using countable-B by (auto simp add: b-def
intro!: range-from-nat-into)
 have open-b: open (b i) for i unfolding b-def using B-balls open-ball from-nat-into[of
B \cup \{\{\}\}\ i by force
 have Union-range-b: \bigcup (range\ b) = -S using B-balls range-b by simp
 {
   fix v r assume ball-in-Compl: ball v r \subseteq -S
   define A where A = f - `ball v r \cap space N
   have dist-less: dist (f x) v < r if x \in A for x using that unfolding A-def
vimage-def by (simp add: dist-commute)
    hence AE-less: AE x \in A in N. norm (f x - v) < r by (auto simp add:
dist-norm)
   have *: A \in sets \ N  unfolding A-def by simp
   have emeasure NA = 0
   proof -
```

assume asm: emeasure N~A>0 hence measure-pos: measure N~A>0 unfolding emeasure-eq-measure by simp

**hence**  $(1 / measure\ N\ A) *_R set-lebesgue-integral\ N\ A\ f - v = (1 / measure\ N\ A) *_R set-lebesgue-integral\ N\ A\ (\lambda x.\ f\ x - v)$  using integrable integrable-const \* by (subst set-integral-diff(2), auto simp add: set-integrable-def set-integral-const[OF\*] algebra-simps intro!: integrable-mult-indicator)

**moreover have** norm  $(\int x \in A. (f \ x - v) \partial N) \leq (\int x \in A. \text{norm } (f \ x - v) \partial N)$  using \* by (auto intro!: integral-norm-bound[of  $N \ \lambda x.$  indicator  $A \ x *_R (f \ x - v), THEN \text{ order-trans}]$  integrable-mult-indicator integrable simp add: set-lebesgue-integral-def)

```
ultimately have norm ((1 \mid measure \mid N \mid A) *_R set-lebesgue-integral \mid N \mid A \mid f
(-v) \le set-lebesgue-integral N A (\lambda x. norm (f x - v)) / measure N A using asm
by (auto intro: divide-right-mono)
      also have ... < set-lebesgue-integral N A (\lambda x. r) / measure N A
        unfolding set-lebesque-integral-def
        using \ asm * integrable \ integrable-const \ AE-less \ measure-pos
      by (intro divide-strict-right-mono integral-less-AE[of - - A] integrable-mult-indicator)
          (fastforce simp add: dist-less dist-norm indicator-def)+
      also have ... = r using * measure-pos by (simp add: set-integral-const)
      finally have dist ((1 / measure N A) *_R set-lebesgue-integral N A f) v < r
by (subst dist-norm)
    hence False using average [OF * measure-pos] by (metis ComplD dist-commute
in-mono mem-ball ball-in-Compl)
     thus ?thesis by fastforce
   qed
 note * = this
   fix b' assume b' \in B
   hence ball-subset-Compl: b' \subseteq -S and ball-radius-pos: \exists v \in D. \exists r > 0. b' =
ball v r using B-balls by (blast, fast)
 \mathbf{note} ** = this
 hence emeasure N (f - b i \cap space N) = 0 for i by (cases b i = \{\}, simp)
(metis UnE singletonD * range-b[THEN eq-refl, THEN range-subsetD])
  hence emeasure N ([]i. f - bi \cap space N) = 0 using open-b by (intro
emeasure-UN-eq-0) fastforce+
 moreover have (\bigcup i. f - b i \cap space N) = f - (\bigcup (range b)) \cap space N by
 ultimately have emeasure N (f - (-S) \cap space N) = 0 using Union-range-b
 hence AEx in N. fx \notin -S using open-Compl[OF assms(2)] by (intro AE-iff-measurable[THEN
iffD2], auto)
 thus ?case by force
qed (simp add: pred-sets2[OF borel-closed] assms(2))
lemma density-zero:
 fixes f::'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology, banach\}
 assumes integrable M f
     and density-0: \bigwedge A. A \in sets \ M \Longrightarrow set-lebesgue-integral M \ A \ f = 0
 shows AE x in M. f x = 0
 using averaging-theorem[OF assms(1), of \{0\}] assms(2)
 by (simp add: scaleR-nonneg-nonneg)
The following lemma shows that densities are unique in Banach spaces.
lemma density-unique-banach:
 fixes f f'::'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology, banach}\}
 assumes integrable M f integrable M f'
```

```
and density-eq: \bigwedge A. A \in sets M \Longrightarrow set-lebesgue-integral M A f = set-lebesgue-integral
M A f'
 shows AE x in M. f x = f' x
proof-
   fix A assume asm: A \in sets M
    hence LINT x|M. indicat-real A x *_R (f x - f' x) = 0 using density-eq
assms(1,2) by (simp add: set-lebesque-integral-def algebra-simps Bochner-Integration.integral-diff[OF]
integrable-mult-indicator(1,1)])
 thus ?thesis using density-zero[OF Bochner-Integration.integrable-diff[OF assms(1,2)]]
by (simp add: set-lebesgue-integral-def)
qed
lemma density-nonneg:
 fixes f::-\Rightarrow b:\{second\-countable\-topology, banach, linorder\-topology, ordered\-real\-vector\}
 assumes integrable M f
     and \bigwedge A. A \in sets \ M \Longrightarrow set-lebesgue-integral M \ A \ f \ge 0
   shows AE x in M. f x > 0
  using averaging-theorem [OF\ assms(1),\ of\ \{0..\},\ OF\ closed-atLeast]\ assms(2)
 by (simp add: scaleR-nonneg-nonneg)
corollary integral-nonneg-eq-0-iff-AE-banach:
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes f[measurable]: integrable M f and nonneg: AE x in M. 0 \le f x
 shows integral<sup>L</sup> M f = 0 \longleftrightarrow (AE x \text{ in } M. f x = 0)
 assume *: integral^L M f = 0
   fix A assume asm: A \in sets M
   have 0 \leq integral^L M (\lambda x. indicator A x *_R f x) using nonneg by (subst inte-
gral-zero[of\ M,\ symmetric],\ intro\ integral-mono-AE-banach\ integrable-mult-indicator
asm f integrable-zero, auto simp add: indicator-def)
  moreover have ... \leq integral^L Mf using nonneg by (intro integral-mono-AE-banach
integrable-mult-indicator asm f, auto simp add: indicator-def)
  ultimately have set-lebesgue-integral MAf = 0 unfolding set-lebesgue-integral-def
using * by force
 }
 thus AE x in M. f x = 0 by (intro density-zero f, blast)
qed (auto simp add: integral-eq-zero-AE)
corollary integral-eq-mono-AE-eq-AE:
  fixes f g :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes integrable M f integrable M g integral L M f = integral L M g AE x in
M. f x < q x
 shows AE x in M. f x = g x
proof -
```

```
define h where h=(\lambda x.\ g\ x-f\ x) have AE\ x in M.\ h\ x=0 unfolding h-def using assms by (subst integral-nonneg-eq-0-iff-AE-banach[symmetric]) auto then show ?thesis unfolding h-def by auto qed end end theory Conditional-Expectation-Banach imports HOL-Probability.Conditional-Expectation\ HOL-Probability.Independent-Family Bochner-Integration-Supplement begin
```

# 4 Conditional Expectation in Banach Spaces

Before we can talk about 'the' conditional expectation, we must define what it means for a function to have a conditional expectation.

```
definition has-cond-exp:: 'a measure \Rightarrow 'a measure \Rightarrow ('a \Rightarrow 'b) \Rightarrow ('a \Rightarrow 'b::{real-normed-vector, second-countable-topology}) \Rightarrow bool where
has-cond-exp M F f g = ((\forall A \in sets F. (\int x \in A. f x \partial M) = (\int x \in A. g x \partial M))

\land integrable M f
\land integrable M g
\land g \in borel-measurable F)
```

This predicate precisely characterizes what it means for a function f to have a conditional expectation g, with respect to the measure M and the sub- $\sigma$ -algebra F.

```
lemma has-cond-expI':

assumes \bigwedge A. A \in sets\ F \Longrightarrow (\int x \in A.\ f\ x\ \partial M) = (\int x \in A.\ g\ x\ \partial M)

integrable M f

integrable M g

g \in borel\text{-}measurable\ F

shows has-cond-exp M F f g

using assms unfolding has-cond-exp-def by simp

lemma has-cond-expD:
assumes has-cond-exp M F f g

shows \bigwedge A. A \in sets\ F \Longrightarrow (\int x \in A.\ f\ x\ \partial M) = (\int x \in A.\ g\ x\ \partial M)

integrable M f

integrable M g

g \in borel\text{-}measurable\ F

using assms unfolding has-cond-exp-def by simp+
```

Now we can use Hilberts  $\epsilon$ -operator to define the conditional expectation,

```
if it exists.
definition cond-exp :: 'a measure \Rightarrow 'a measure \Rightarrow ('a \Rightarrow 'b) \Rightarrow ('a \Rightarrow 'b::{banach,
second-countable-topology}) where
 cond-exp M F f = (if \exists q. has\text{-}cond\text{-}exp M F f q then (SOME q. has\text{-}cond\text{-}exp M
F f g) else (\lambda -. \theta))
lemma borel-measurable-cond-exp[measurable]: cond-exp M F f \in borel-measurable
 \mathbf{by}\ (\mathit{metis}\ \mathit{cond-exp-def}\ \mathit{someI}\ \mathit{has-cond-exp-def}\ \mathit{borel-measurable-const})
lemma integrable-cond-exp[intro]: integrable\ M\ (cond-exp\ M\ F\ f)
 by (metis\ cond\text{-}exp\text{-}def\ has\text{-}cond\text{-}expD(3)\ integrable\text{-}zero\ some I)
lemma set-integrable-cond-exp[intro]:
  assumes A \in sets M
  shows set-integrable M A (cond-exp M F f) using integrable-mult-indicator OF
assms integrable-cond-exp, of F f] by (auto simp add: set-integrable-def intro!: in-
tegrable-mult-indicator[OF\ assms\ integrable-cond-exp])
lemma has-cond-exp-self:
 assumes integrable M f
 shows has-cond-exp M (vimage-algebra (space M) f borel) f f
 using assms by (auto intro!: has-cond-expI' measurable-vimage-algebra1)
\mathbf{lemma}\ \mathit{has\text{-}cond\text{-}exp\text{-}sets\text{-}cong}\text{:}
 assumes sets F = sets G
 shows has-cond-exp M F = has-cond-exp M G
 using assms unfolding has-cond-exp-def by force
lemma cond-exp-sets-cong:
  assumes sets F = sets G
 shows AE x in M. cond-exp M F f x = cond-exp M G f x
  by (intro AE-I2, simp add: cond-exp-def has-cond-exp-sets-cong[OF assms, of
M
{\bf context}\ sigma-finite-subalgebra
begin
lemma borel-measurable-cond-exp'[measurable]: cond-exp M F f \in borel-measurable
 by (metis cond-exp-def some I has-cond-exp-def borel-measurable-const subalg mea-
surable-from-subalg)
lemma cond-exp-null:
 assumes \nexists g. has-cond-exp M F f g
 shows cond-exp M F f = (\lambda - 0)
 unfolding cond-exp-def using assms by argo
```

We state the tower property of the conditional expectation in terms of the

```
predicate has-cond-exp.
lemma has-cond-exp-nested-subalg:
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
 assumes subalgebra G F has-cond-exp M F f h has-cond-exp M G f h'
 shows has-cond-exp M F h' h
 by (intro has-cond-expI') (metis assms has-cond-expD in-mono subalgebra-def)+
The following lemma shows that the conditional expectation is unique as an
element of L1, given that it exists.
lemma has-cond-exp-charact:
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
 assumes has-cond-exp M F f g
 shows has\text{-}cond\text{-}exp\ M\ F\ f\ (cond\text{-}exp\ M\ F\ f)
       AE \ x \ in \ M. \ cond\text{-}exp \ M \ F \ f \ x = g \ x
proof -
  show cond-exp: has-cond-exp M F f (cond-exp M F f) using assms someI
cond-exp-def by metis
 let ?MF = restr-to-subalg\ M\ F
 interpret sigma-finite-measure ?MF by (rule sigma-fin-subalg)
   fix A assume A \in sets ?MF
    then have [measurable]: A \in sets \ F \ using \ sets-restr-to-subalg[OF \ subalg] by
simp
   have (\int x \in A. \ g \ x \ \partial ?MF) = (\int x \in A. \ g \ x \ \partial M) using assms subalg by (auto
simp add: integral-subalgebra2 set-lebesgue-integral-def dest!: has-cond-expD)
    also have ... = (\int x \in A. \ cond\text{-}exp \ M \ F \ f \ x \ \partial M) using assms cond-exp by
(simp add: has-cond-exp-def)
   also have ... = (\int x \in A. \ cond\text{-}exp \ M F \ f \ x \ \partial ?MF) using subalg by (auto simp
add: integral-subalgebra2 set-lebesgue-integral-def)
    finally have (\int x \in A. \ g \ x \ \partial ?MF) = (\int x \in A. \ cond\text{-exp} \ M \ F \ f \ x \ \partial ?MF) by
simp
 hence AE \ x \ in \ ?MF. \ cond-exp \ M \ F \ f \ x = g \ x \ using \ cond-exp \ assms \ subalg \ by
(intro density-unique-banach, auto dest: has-cond-expD intro!: integrable-in-subalg)
  then show AE \times in M. cond-exp M F f \times g \times using AE-restr-to-subalg[OF]
subalg] by simp
qed
corollary cond-exp-charact:
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
 assumes \bigwedge A. A \in sets \ F \Longrightarrow (\int x \in A. \ f \ x \ \partial M) = (\int x \in A. \ g \ x \ \partial M)
         integrable M f
         integrable M g
         g \in borel-measurable F
   shows AE x in M. cond-exp M F f x = g x
  by (intro has-cond-exp-charact has-cond-expI' assms) auto
```

Identity on F-measurable functions:

```
If an integrable function f is already F-measurable, then cond-exp M F f =
f t-a.e. This is a corollary of the lemma on the characterization of cond-exp.
corollary cond-exp-F-meas[intro, simp]:
 fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
 assumes integrable M f
        f \in borel-measurable F
   shows AE x in M. cond-exp M F f x = f x
 by (rule cond-exp-charact, auto intro: assms)
Congruence
lemma has-cond-exp-cong:
 assumes integrable M f \land x. x \in space M \Longrightarrow f x = g x has-cond-exp M F g h
 shows has\text{-}cond\text{-}exp\ M\ F\ f\ h
proof (intro\ has\text{-}cond\text{-}expI'[OF\ -\ assms(1)])
 fix A assume asm: A \in sets F
 hence set-lebesque-integral MAf = set-lebesque-integral MAg by (intro set-lebesque-integral-cong)
(meson assms(2) subalq in-mono subalqebra-def sets.sets-into-space subalqebra-def
subsetD)+
 thus set-lebesque-integral M A f = set-lebesque-integral M A h using asm assms(3)
by (simp add: has-cond-exp-def)
qed (auto simp add: has-cond-expD[OF assms(3)])
lemma cond-exp-cong:
 fixes f :: 'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology,banach}\}
 assumes integrable M f integrable M g \bigwedge x. x \in space M \Longrightarrow f x = g x
 shows AE x in M. cond-exp M F f x = cond-exp M F g x
proof (cases \exists h. has-cond-exp M F f h)
 case True
  then obtain h where h: has-cond-exp M F f h has-cond-exp M F g h using
has-cond-exp-cong assms by metis
 show ?thesis using h[THEN\ has-cond-exp-charact(2)] by fastforce
next
 case False
 moreover have \nexists h. has-cond-exp M F g h using False has-cond-exp-cong assms
 ultimately show ?thesis unfolding cond-exp-def by auto
qed
lemma has-cond-exp-conq-AE:
 assumes integrable M f AE x in M. f x = g x has-cond-exp M F g h
 shows has\text{-}cond\text{-}exp\ M\ F\ f\ h
 using assms(1,2) subalg subalgebra-def subset-iff
 by (intro has-cond-expI', subst set-lebesque-integral-cong-AE[OF-assms(1)]THEN
borel-measurable-integrable|\ borel-measurable-integrable(1)[OF\ has-cond-expD(2)]OF
assms(3)]]])
   (fast\ intro:\ has\text{-}cond\text{-}expD[OF\ assms(3)]\ integrable\text{-}cong\text{-}AE\text{-}imp[OF\ -\ -\ AE\text{-}symmetric]})+
lemma has-cond-exp-cong-AE':
```

assumes  $h \in borel$ -measurable  $F \land AE \ x \ in \ M. \ h \ x = h' \ x \ has$ -cond-exp  $M \ F \ f \ h'$ 

```
shows has\text{-}cond\text{-}exp\ M\ F\ f\ h
  using assms(1, 2) subalg subalgebra-def subset-iff
 using AE-restr-to-subalg2[OF subalg assms(2)] measurable-from-subalg
 by (intro has-cond-expI', subst set-lebesque-integral-cong-AE[OF - measurable-from-subalg(1,1)[OF
subalg], OF - assms(1) \ has-cond-expD(4)[OF \ assms(3)]])
   (fast\ intro:\ has-cond-expD[OF\ assms(3)]\ integrable-cong-AE-imp[OF\ -\ -\ AE-symmetric])+
lemma cond-exp-cong-AE:
  fixes f :: 'a \Rightarrow 'b :: \{ second\text{-}countable\text{-}topology, banach \}
 assumes integrable M f integrable M g AE x in M. f x = g x
 shows AE x in M. cond-exp M F f x = cond-exp M F g x
proof (cases \exists h. has\text{-}cond\text{-}exp M F f h)
 case True
  then obtain h where h: has-cond-exp M F f h has-cond-exp M F g h using
has-cond-exp-conq-AE assms by (metis (mono-tags, lifting) eventually-mono)
 show ?thesis using h[THEN\ has-cond-exp-charact(2)] by fastforce
next
 {f case}\ {\it False}
  moreover have \nexists h. has-cond-exp M F g h using False has-cond-exp-cong-AE
assms by auto
  ultimately show ?thesis unfolding cond-exp-def by auto
qed
The conditional expectation operator on the reals, real-cond-exp, satisfies
the conditions of the conditional expectation as we have defined it.
lemma has-cond-exp-real:
 fixes f :: 'a \Rightarrow real
 assumes integrable M f
 shows has\text{-}cond\text{-}exp\ M\ F\ f\ (real\text{-}cond\text{-}exp\ M\ F\ f)
 by (intro has-cond-expI', auto intro!: real-cond-exp-intA assms)
lemma \ cond-exp-real[intro]:
 fixes f :: 'a \Rightarrow real
 assumes integrable M f
 shows AE \ x \ in \ M. \ cond-exp \ M \ F \ f \ x = real-cond-exp \ M \ F \ f \ x
 using has-cond-exp-charact has-cond-exp-real assms by blast
lemma cond-exp-cmult:
  fixes f :: 'a \Rightarrow real
 assumes integrable M f
 shows AE \ x \ in \ M. \ cond\text{-}exp \ M \ F \ (\lambda x. \ c * f \ x) \ x = c * cond\text{-}exp \ M \ F \ f \ x
  using real-cond-exp-cmult[OF\ assms(1),\ of\ c]\ assms(1)[THEN\ cond-exp-real]
assms(1)[THEN integrable-mult-right, THEN cond-exp-real, of c] by fastforce
```

#### 4.1 Existence

Showing the existence is a bit involved. Specifically, what we aim to show is that  $has\text{-}cond\text{-}exp\ M\ F\ f\ (cond\text{-}exp\ M\ F\ f)$  holds for any Bochner-integrable f. We will employ the standard machinery of measure theory. First, we will

prove existence for indicator functions. Then we will extend our proof by linearity to simple functions. Finally we use a limiting argument to show that the conditional expectation exists for all Bochner-integrable functions.

#### Indicator functions

```
lemma has-cond-exp-indicator:
  assumes A \in sets \ M \ emeasure \ M \ A < \infty
  shows has-cond-exp M F (\lambda x. indicat-real A x *_R y) (\lambda x. real-cond-exp M F
(indicator\ A)\ x *_R y)
proof (intro has-cond-expI', goal-cases)
  case (1 B)
  have \int x \in B. (indicat-real A \times R y) \partial M = (\int x \in B. indicat-real A \times \partial M) *_R
y using assms by (intro set-integral-scaleR-left, meson 1 in-mono subalg subalge-
bra-def, blast)
  also have ... = (\int x \in B. \ real\text{-}cond\text{-}exp \ M \ F \ (indicator \ A) \ x \ \partial M) *_R y \ using 1
assms by (subst real-cond-exp-intA, auto)
  also have ... = \int x \in B. (real-cond-exp M F (indicator A) x *_R y) \partial M using
assms by (intro set-integral-scaleR-left[symmetric], meson 1 in-mono subalg subal-
gebra-def, blast)
  finally show ?case.
next
  case 2
  show ?case using integrable-scaleR-left integrable-real-indicator assms by blast
next
  case \beta
 show ?case using assms by (intro integrable-scaleR-left, intro real-cond-exp-int,
blast+)
\mathbf{next}
  case 4
 show ?case by (intro borel-measurable-scaleR, intro Conditional-Expectation.borel-measurable-cond-exp,
simp)
qed
lemma cond-exp-indicator[intro]:
  fixes y :: 'b::{second-countable-topology,banach}
  assumes [measurable]: A \in sets \ M \ emeasure \ M \ A < \infty
 shows AE \ x \ in \ M. \ cond-exp \ M \ F \ (\lambda x. \ indicat-real \ A \ x *_R \ y) \ x = cond-exp \ M \ F
(indicator\ A)\ x*_R\ y
proof -
 have AE x in M. cond-exp MF (\lambda x. indicat-real A x *_R y) x = real-cond-exp MF
(indicator\ A)\ x*_R\ y\ {\bf using}\ has\text{-}cond\text{-}exp\text{-}indicator[OF\ assms]\ has\text{-}cond\text{-}exp\text{-}charact
by blast
 thus ?thesis using cond-exp-real[OF integrable-real-indicator, OF assms] by fast-
force
qed
Addition
\mathbf{lemma}\ \mathit{has}\text{-}\mathit{cond}\text{-}\mathit{exp}\text{-}\mathit{add}\text{:}
```

fixes  $fg:: 'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology,banach}\}$ 

```
assumes has-cond-exp M F f f' has-cond-exp M F g g'
 shows has-cond-exp M F (\lambda x. f x + g x) (\lambda x. f' x + g' x)
proof (intro has-cond-expI', goal-cases)
  case (1 A)
  have \int x \in A. (f x + g x) \partial M = (\int x \in A \cdot f x \partial M) + (\int x \in A \cdot g x \partial M) using
assms[THEN\ has-cond-expD(2)]\ subalg\ 1\ \mathbf{by}\ (intro\ set-integral-add(2),\ auto\ simp)
add: subalgebra-def set-integrable-def intro: integrable-mult-indicator)
  also have ... = (\int x \in A. \ f' \ x \ \partial M) + (\int x \in A. \ g' \ x \ \partial M) using assms[THEN]
has\text{-}cond\text{-}expD(1)[OF - 1]] by argo
 also have ... = \int x \in A. (f'x + g'x)\partial M using assms[THEN\ has-cond-expD(3)]
subalg 1 by (intro set-integral-add(2)[symmetric], auto simp add: subalgebra-def
set-integrable-def intro: integrable-mult-indicator)
 finally show ?case.
next
 case 2
 show ?case by (metis Bochner-Integration.integrable-add assms has-cond-expD(2))
next
 case 3
 show ?case by (metis Bochner-Integration.integrable-add assms has-cond-expD(3))
next
 case 4
 show ?case using assms borel-measurable-add has-cond-expD(4) by blast
qed
lemma has-cond-exp-scaleR-right:
  fixes f :: 'a \Rightarrow 'b :: \{ second\text{-}countable\text{-}topology, banach \}
 assumes has-cond-exp M F f f'
 shows has-cond-exp M F (\lambda x. c *_R f x) (\lambda x. c *_R f' x)
 using has-cond-expD[OF assms] by (intro has-cond-expI', auto)
lemma cond-exp-scaleR-right:
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
 assumes integrable M f
 shows AE \ x \ in \ M. \ cond-exp \ M \ F \ (\lambda x. \ c *_R f \ x) \ x = c *_R \ cond-exp \ M \ F \ f \ x
proof (cases \exists f'. has-cond-exp M F f f')
 case True
 then show ?thesis using assms has-cond-exp-charact has-cond-exp-scaleR-right
by metis
\mathbf{next}
 case False
 show ?thesis
 proof (cases \ c = \theta)
   case True
   then show ?thesis by simp
  next
   case c-nonzero: False
   have \nexists f'. has-cond-exp M F (\lambda x. c *_R f x) f'
   proof (standard, goal-cases)
     case 1
```

```
then obtain f' where f': has-cond-exp M F (\lambda x. c *_R f x) f' by blast have has-cond-exp M F f (\lambda x. inverse c *_R f' x) using has-cond-expD[OF f'] divideR-right[OF c-nonzero] assms by (intro has-cond-expI', auto) then show ?case using False by blast qed then show ?thesis using cond-exp-null[OF False] cond-exp-null by force qed qed lemma cond-exp-uminus: fixes f :: 'a \Rightarrow 'b::\{second-countable-topology,banach\} assumes integrable M f shows AE x in M. cond-exp M F (\lambda x. -f x) x = - cond-exp M F f x using cond-exp-scaleR-right[OF assms, of -1] by force
```

Together with the induction scheme *integrable-simple-function-induct*, we can show that the conditional expectation of an integrable simple function exists.

```
corollary has-cond-exp-simple:
fixes f:: 'a \Rightarrow 'b:: \{second-countable-topology, banach\}
assumes simple-function\ M\ f\ emeasure\ M\ \{y \in space\ M.\ f\ y \neq 0\} \neq \infty
shows has-cond-exp\ M\ F\ f\ (cond-exp\ M\ F\ f)
using assms
proof (induction rule: integrable-simple-function-induct)
case (cong\ f\ g)
then show ?case\ using\ has-cond-exp-cong\ by\ (metis\ (no-types,\ opaque-lifting)
Bochner-Integration.integrable-cong\ has-cond-expD(2)\ has-cond-exp-charact(1))
next
case (indicator\ A\ y)
then show ?case\ using\ has-cond-exp-charact[OF\ has-cond-exp-indicator]\ by\ fast
next
case (add\ u\ v)
then show ?case\ using\ has-cond-exp-add\ has-cond-exp-charact(1)\ by\ blast
qed
```

Now comes the most difficult part. Given a convergent sequence of integrable simple functions s, we must show that the sequence  $\lambda n$ . cond-exp M F (s n) is also convergent. Furthermore, we must show that this limit satisfies the properties of a conditional expectation. Unfortunately, we will only be able to show that this sequence convergences in the L1-norm. Luckily, this is enough to show that the operator cond-exp M F preserves limits as a function from L1 to L1.

In anticipation of this result, we show that the conditional expectation operator is a contraction for simple functions. We first reformulate the lemma real-cond-exp-abs, which shows the statement for real-valued functions, using our definitions. Then we show the statement for simple functions via induction.

```
assumes integrable[measurable]: integrable M f
   shows AE \times in M. norm (cond-exp M F f x) \leq cond-exp M F (<math>\lambda x. norm (f x)) x \in AE \times in M.
proof-
   have int: integrable M (\lambda x. norm (f x)) using assms by blast
  have *: AE x in M. 0 \le cond\text{-}exp M F (\lambda x. norm (f x)) x using cond\text{-}exp\text{-}real[THEN]
AE-symmetric, OF integrable-norm[OF integrable] real-cond-exp-qe-c[OF integrable-norm[OF
integrable, of 0] norm-ge-zero by fastforce
    have **: A \in sets \ F \Longrightarrow \int x \in A. |f \ x| \ \partial M = \int x \in A. real-cond-exp M \ F \ (\lambda x).
norm (f x)) x \partial M for A unfolding real-norm-def using assms integrable-abs
real-cond-exp-intA by blast
   have norm-int: A \in sets \ F \Longrightarrow (\int x \in A. \ |f \ x| \ \partial M) = (\int x \in A. \ |f \ x| \ \partial M) for A
using assms by (intro nn-set-integral-eq-set-integral[symmetric], blast, fastforce)
(meson\ subalq\ subalqebra-def\ subset D)
  have AEx in M. real-cond-exp MF(\lambda x. norm(fx)) x \geq 0 using int real-cond-exp-ge-c
by force
  hence cond-exp-norm-int: A \in sets \ F \Longrightarrow (\int x \in A. \ real\text{-}cond\text{-}exp \ M \ F \ (\lambda x. \ norm
(f x) (f x)
assms by (intro nn-set-integral-eq-set-integral[symmetric], blast, fastforce) (meson
subalg\ subalgebra-def\ subset D)
    have A \in sets \ F \Longrightarrow \int x \in A. |f x| \partial M = \int x \in A. real-cond-exp M F (\lambda x).
norm\ (f\ x))\ x\ \partial M for A using ** norm\text{-}int\ cond\text{-}exp\text{-}norm\text{-}int by (auto simp
add: nn-integral-set-ennreal)
   moreover have (\lambda x. \ ennreal \ |f \ x|) \in borel-measurable \ M by measurable
  moreover have (\lambda x. \ ennreal \ (real\text{-}cond\text{-}exp \ M \ F \ (\lambda x. \ norm \ (f \ x)) \ x)) \in borel\text{-}measurable
F by measurable
  ultimately have AE x in M. nn-cond-exp MF (\lambda x. ennreal |fx|) x = real-cond-exp
M F (\lambda x. norm (f x)) x by (intro nn-cond-exp-charact[THEN AE-symmetric],
auto)
   hence AE \times in M. nn-cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-exp M F (\lambda x. ennreal | f x |) x \leq cond-
norm (f x)) x using cond-exp-real [OF int] by force
   moreover have AE x in M. |real-cond-exp M F f x| = norm (cond-exp M F f x)
unfolding real-norm-def using cond-exp-real[OF assms] * by force
   (\lambda x.\ norm\ (fx))\ x\ using\ real-cond-exp-abs[OF\ assms[THEN\ borel-measurable-integrable]]
by fastforce
   hence AE \ x \ in \ M. \ enn2real \ (ennreal \ (norm \ (cond-exp \ M \ F \ f \ x))) \le enn2real
(cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (f\ x))\ x) using ennreal-le-iff2 by force
   thus ?thesis using * by fastforce
qed
lemma cond-exp-contraction-simple:
   fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
   assumes simple-function M f emeasure M \{y \in space M. f y \neq 0\} \neq \infty
   shows AE \ x \ in \ M. \ norm \ (cond-exp \ M \ F \ f \ x) \le cond-exp \ M \ F \ (\lambda x. \ norm \ (f \ x)) \ x
```

**lemma** cond-exp-contraction-real:

fixes  $f :: 'a \Rightarrow real$ 

```
proof (induction rule: integrable-simple-function-induct)
  case (cong f g)
  hence ae: AE x in M. f x = g x by blast
 hence AEx in M, cond-exp MFfx = cond-exp MFqx using cond has-cond-exp-simple
by (subst cond-exp-cong-AE) (auto intro!: has-cond-expD(2))
  hence AE \ x \ in \ M. \ norm \ (cond-exp \ M \ F \ f \ x) = norm \ (cond-exp \ M \ F \ g \ x) by
force
  moreover have AE x in M. cond-exp M F (\lambda x. norm (f x)) x = cond-exp M F
(\lambda x.\ norm\ (g\ x))\ x using ae cong has-cond-exp-simple by (subst cond-exp-cong-AE)
(auto\ dest:\ has-cond-expD)
  ultimately show ? case using cong(6) by fastforce
next
  case (indicator\ A\ y)
  hence AE \ x \ in \ M. \ cond-exp \ M \ F \ (\lambda a. \ indicator \ A \ a *_R \ y) \ x = cond-exp \ M \ F
(indicator A) x *_{B} y by blast
 hence *: AEx in M. norm (cond-exp MF (\lambda a. indicat-real A \ a *_R y) x) \leq norm y
* cond-exp\ M\ F\ (\lambda x.\ norm\ (indicat\text{-}real\ A\ x))\ x\ \mathbf{using}\ cond\text{-}exp\text{-}contraction\text{-}real[OF]
integrable-real-indicator, OF indicator by fastforce
 have AE \times in M. norm \ y \times cond-exp \ MF \ (\lambda x. \ norm \ (indicat-real \ A \times)) \ x = norm
y * real\text{-}cond\text{-}exp \ M \ F \ (\lambda x. \ norm \ (indicat\text{-}real \ A \ x)) \ x \ using \ cond\text{-}exp\text{-}real[OF]
integrable-real-indicator, OF indicator by fastforce
  moreover have AE x in M. cond-exp M F (\lambda x. norm y * norm (indicat-real
(A \ x)) x = real\text{-}cond\text{-}exp \ M \ F \ (\lambda x. \ norm \ y * norm \ (indicat\text{-}real \ A \ x)) \ x \ using
indicator by (intro cond-exp-real, auto)
 ultimately have AE \times in M. norm y * cond-exp M F (\lambda x. norm (indicat-real A))
x)) x = cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ y*norm\ (indicat\text{-}real\ A\ x))\ x\ using\ real\text{-}cond\text{-}exp\text{-}cmult[of
\lambda x. \ norm \ (indicat-real \ A \ x) \ norm \ y] \ indicator \ \mathbf{by} \ fastforce
 moreover have (\lambda x. norm \ y * norm \ (indicat-real \ A \ x)) = (\lambda x. norm \ (indicat-real \ A \ x))
A x *_R y) by force
  ultimately show ?case using * by force
next
  case (add\ u\ v)
 have AE \times in M. norm (cond-exp M F (\lambda a. u a + v a) \times) = norm (cond-exp M F (\lambda a. u a + v a) \times)
Fux + cond-exp MFvx) using has-cond-exp-charact(2)[OF has-cond-exp-add,
OF has-cond-exp-simple (1,1), OF add (1,2,3,4) by fastforce
  moreover have AE \ x \ in \ M. \ norm \ (cond-exp \ M \ F \ u \ x + cond-exp \ M \ F \ v \ x) <
norm (cond\text{-}exp \ M \ F \ u \ x) + norm (cond\text{-}exp \ M \ F \ v \ x) using norm\text{-}triangle\text{-}ineq
by blast
 \mathbf{moreover} \ \mathbf{have} \ AE \ x \ in \ M. \ norm \ (\mathit{cond-exp} \ M \ F \ u \ x) + \mathit{norm} \ (\mathit{cond-exp} \ M \ F \ v
x > cond-exp M F (\lambda x. norm (u x)) x + cond-exp M F (\lambda x. norm (v x)) x using
add(6,7) by fastforce
 moreover have AE x in M. cond-exp M F (\lambda x. norm (u x)) x + cond-exp M F
(\lambda x. \ norm \ (v \ x)) \ x = cond\text{-}exp \ M \ F \ (\lambda x. \ norm \ (u \ x) + norm \ (v \ x)) \ x \ using \ in-
tegrable-simple-function [OF add(1,2)] integrable-simple-function [OF add(3,4)] by
(intro\ has\text{-}cond\text{-}exp\text{-}charact(2)[OF\ has\text{-}cond\text{-}exp\text{-}add[OF\ has\text{-}cond\text{-}exp\text{-}charact(1,1)],}
THEN AE-symmetric, auto intro: has-cond-exp-real)
 moreover have AE x in M. cond-exp MF (\lambda x. norm (u x) + norm (v x)) x =
```

using assms

```
cond-exp\ M\ F\ (\lambda x.\ norm\ (u\ x+v\ x))\ x\ using\ add(5)\ integrable-simple-function[OF]
add(1,2)] integrable-simple-function[OF add(3,4)] by (intro cond-exp-cong, auto)
 ultimately show ?case by force
qed
```

 $\mathbf{lemma}\ \mathit{has\text{-}cond\text{-}exp\text{-}simple\text{-}lim} \colon$ 

```
fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
assumes integrable[measurable]: integrable M f
    and \bigwedge i. simple-function M (s i)
    and \bigwedge i. emeasure M \{ y \in space M. \ s \ i \ y \neq 0 \} \neq \infty
    and \bigwedge x. \ x \in space \ M \Longrightarrow (\lambda i. \ s \ i \ x) \longrightarrow f \ x
    and \bigwedge x \ i. \ x \in space \ M \Longrightarrow norm \ (s \ i \ x) \le 2 * norm \ (f \ x)
obtains r
where strict-mono r has-cond-exp M F f (\lambda x. lim (\lambda i. cond-exp M F (s (r i))
```

AE x in M. convergent ( $\lambda i$ . cond-exp M F (s (r i)) x)

proof -

have [measurable]:  $(s\ i) \in borel-measurable\ M$  for  $i\ using\ assms(2)\ by\ (simp\ partial or i)$ add: borel-measurable-simple-function)

have integrable-s: integrable  $M(\lambda x. six)$  for i using assms integrable-simple-function **bv** blast

have integrable-4f: integrable M ( $\lambda x$ . 4 \* norm (f x)) using assms(1) by simp have integrable-2f: integrable M ( $\lambda x$ . 2 \* norm (f x)) using assms(1) by simp have integrable-2-cond-exp-norm-f: integrable M ( $\lambda x$ . 2 \* cond-exp M F ( $\lambda x$ . norm (f x)) x) by fast

have emeasure  $M \{ y \in space M. \ s \ i \ y - s \ j \ y \neq 0 \} \leq emeasure M \{ y \in space M \}$ space M. s i  $y \neq 0$ } + emeasure M { $y \in space M. s j y \neq 0$ } for i j using  $simple-function D(2)[OF\ assms(2)]$  by  $(intro\ order-trans[OF\ emeasure-mono\ emea$ sure-subadditive, auto)

hence fin-sup: emeasure M  $\{y \in space M. \ s \ i \ y - s \ j \ y \neq 0\} \neq \infty$  for i j using assms(3) by  $(metis\ (mono-tags)\ ennreal-add-eq-top\ linorder-not-less$ top.not-eq-extremum infinity-ennreal-def)

have emeasure M  $\{y \in space M. norm (s i y - s j y) \neq 0\} \leq emeasure M$  $\{y \in space \ M. \ s \ i \ y \neq 0\} + emeasure \ M \ \{y \in space \ M. \ s \ j \ y \neq 0\} \$ for  $i \ j \$ using  $simple-functionD(2)[OF\ assms(2)]$  by (intro\ order-trans[OF\ emeasure-mono emeasure-subadditive, auto)

hence fin-sup-norm: emeasure M  $\{y \in space M. norm (s i y - s j y) \neq 0\} \neq \infty$ for i j using assms(3) by (metis (mono-tags) ennreal-add-eq-top linorder-not-less top.not-eq-extremum infinity-ennreal-def)

have Cauchy: Cauchy  $(\lambda n. \ s \ n \ x)$  if  $x \in space \ M$  for x using assms(4) LIM-SEQ-imp-Cauchy that by blast

hence bounded-range-s: bounded (range  $(\lambda n. \ s \ n \ x)$ ) if  $x \in space \ M$  for x using that cauchy-imp-bounded by fast

Since the sequence  $\lambda n$ . s n x is Cauchy for almost all x, we know that the diameter tends to zero almost everywhere.

Dominated convergence tells us that the integral of the diameter also converges to zero.

```
have AE x in M. (\lambda n. diameter \{s \ i \ x \mid i. \ n \leq i\}) \longrightarrow \theta using Cauchy
cauchy-iff-diameter-tends-to-zero-and-bounded by fast
  moreover have (\lambda x. \ diameter \ \{s \ i \ x \ | i. \ n \leq i\}) \in borel-measurable M \ for n
using bounded-range-s borel-measurable-diameter by measurable
 moreover have AE x in M. norm (diameter \{s \mid i \mid i \mid n \leq i\}) \leq 4 * norm (f
x) for n
 proof -
     fix x assume x: x \in space M
       have diameter \{s \ i \ x \ | i. \ n \leq i\} \leq 2 * norm (f \ x) + 2 * norm (f \ x)
by (intro diameter-le, blast, subst dist-norm[symmetric], intro dist-triangle3[THEN
order-trans, of 0, intro add-mono) (auto intro: assms(5)[OF x])
      hence norm (diameter \{s \ i \ x \ | i. \ n \leq i\}) \leq 4 * norm (f \ x) using diame-
ter-ge-0[OF\ bounded-subset[OF\ bounded-range-s],\ OF\ x,\ of\ \{s\ i\ x\ | i.\ n\leq i\}] by
force
   }
   thus ?thesis by fast
 aed
 ultimately have diameter-tendsto-zero: (\lambda n. \ LINT \ x | M. \ diameter \ \{s \ i \ x \mid i. \ n \le a \})
i\}) \longrightarrow 0 by (intro integral-dominated-convergence OF borel-measurable-const of
0] - integrable-4f, simplified]) (fast+)
```

have diameter-integrable: integrable M ( $\lambda x$ . diameter  $\{s \ i \ x \mid i. \ n \leq i\}$ ) for n using assms(1,5)

 $\mathbf{by}\ (intro\ integrable-bound-diameter[OF\ bounded-range-s\ integrable-2f],\ auto)$ 

have dist-integrable: integrable M ( $\lambda x$ . dist (s i x) (s j x)) for i j using assms(5) dist-triangle3 [of s i - - 0, THEN order-trans, OF add-mono, of - 2 \* norm (f -)] by (intro Bochner-Integration.integrable-bound[OF integrable-4f]) fastforce+

Since cond-exp M F is a contraction for simple functions, the following sequence of integral values is also Cauchy.

This follows, since the distance between the terms of this sequence are always less than or equal to the diameter, which itself converges to zero.

Hence, we obtain a subsequence which is Cauchy almost everywhere.

```
hence diameter \{s \ i \ x \mid i.\ N \leq i\} = (SUP\ (i,j) \in \{N..\} \times \{N..\}.\ dist\ (s\ i \ x)\ (s\ j\ x)) unfolding diameter-def by auto
```

moreover have  $(SUP\ (i,j) \in \{N..\} \times \{N..\}.\ dist\ (s\ i\ x)\ (s\ j\ x)) \ge dist\ (s\ i\ x)\ (s\ j\ x)$  using asm bounded-imp-bdd-above [OF bounded-imp-dist-bounded, OF bounded-range-s] by (intro\ cSup-upper, auto)

ultimately have diameter  $\{s \ i \ x \mid i.\ N \leq i\} \geq dist\ (s\ i \ x)\ (s\ j \ x)$  by presburger  $\}$ 

hence LINT x|M. dist  $(s\ i\ x)\ (s\ j\ x) < e\ \text{if}\ i \geq N\ j \geq N\ \text{for}\ i\ j\ \text{using}$  that  $*\ \text{by}\ (intro\ integral-mono}[OF\ dist-integrable\ diameter-integrable,\ THEN\ order.strict-trans1],\ blast+)$ 

moreover have LINT x|M. norm (cond-exp M F (s i) x – cond-exp M F (s j) x)  $\leq$  LINT x|M. dist (s i x) (s j x) for i j proof –

have LINT x|M. norm (cond-exp M F (s i) x – cond-exp M F (s j) x) = LINT x|M. norm (cond-exp M F (s i) x + – 1  $*_R$  cond-exp M F (s j) x) unfolding dist-norm by simp

also have ... = LINT x|M. norm (cond-exp M F ( $\lambda x$ . s i x - s j x) x) using has-cond-exp-charact(2)[OF has-cond-exp-add[OF - has-cond-exp-scaleR-right, OF has-cond-exp-charact(1,1), OF has-cond-exp-simple(1,1)[OF assms(2,3)]], THEN AE-symmetric, of i -1 j] by (intro integral-cong-AE) force+

also have ...  $\leq LINT \ x | M$ .  $cond-exp \ M \ F \ (\lambda x. \ norm \ (s \ i \ x - s \ j \ x)) \ x \ using <math>cond-exp\text{-}contraction\text{-}simple[OF - fin\text{-}sup, of i \ j] } \ integrable\text{-}cond\text{-}exp \ assms(2) \ by \ (intro \ integral\text{-}mono\text{-}AE, \ fast+)$ 

also have ... =  $LINT \ x | M$ .  $norm \ (s \ i \ x - s \ j \ x)$  unfolding set-integral-space(1)[OF integrable-cond-exp, symmetric] set-integral-space[OF dist-integrable[unfolded dist-norm], symmetric] by (intro has-cond-expD(1)[OF has-cond-exp-simple[OF - fin-sup-norm], symmetric]) (metis assms(2) simple-function-compose1 simple-function-diff, metis  $sets.top\ subalq\ subalqebra-def$ )

finally show ?thesis unfolding dist-norm.

qed

ultimately show ?thesis using order.strict-trans1 by meson qed

then obtain r where strict-mono-r: strict-mono r and AE-Cauchy: AE x in M. Cauchy ( $\lambda i$ . cond-exp M F (s (r i)) x)

by (rule cauchy-L1-AE-cauchy-subseq[OF integrable-cond-exp], auto)

hence ae-lim-cond-exp:  $AE \ x \ in \ M$ .  $(\lambda n. \ cond-exp \ M \ F \ (s \ (r \ n)) \ x) \longrightarrow lim \ (\lambda n. \ cond-exp \ M \ F \ (s \ (r \ n)) \ x)$  using Cauchy-convergent-iff convergent-LIMSEQ-iff by fastforce

Now that we have a candidate for the conditional expectation, we must show that it actually has the required properties.

Dominated convergence shows that this limit is indeed integrable.

Here, we again use the fact that conditional expectation is a contraction on simple functions.

**have** cond-exp-bounded: AE x in M. norm (cond-exp M F (s (r n)) x)  $\leq$  cond-exp M F ( $\lambda x$ . 2 \* norm (f x)) x **for** n

```
proof -
      have AE \ x \ in \ M. \ norm \ (cond-exp \ M \ F \ (s \ (r \ n)) \ x) \leq cond-exp \ M \ F \ (\lambda x. \ norm
(s(r n) x) x by (rule\ cond\text{-}exp\text{-}contraction\text{-}simple[OF\ assms(2,3)])
       moreover have AE x in M. real-cond-exp M F (\lambda x. norm (s (r n) x)) x \le
real-cond-exp\ M\ F\ (\lambda x.\ 2*norm\ (f\ x))\ x\ {\bf using}\ integrable-s\ integrable-2f\ assms(5)
by (intro real-cond-exp-mono, auto)
        ultimately show ?thesis using cond-exp-real[OF integrable-norm, OF inte-
grable-s, of r n cond-exp-real OF integrable-2f by force
   qed
    have lim-integrable: integrable M (\lambda x. lim (\lambda i. cond-exp M F (s (r i)) x))
by (intro integrable-dominated-convergence OF - borel-measurable-cond-exp' inte-
grable-cond-exp ae-lim-cond-exp cond-exp-bounded, simp)
Moreover, we can use the DCT twice to show that the conditional expecta-
tion property holds, i.e. the value of the integral of the candidate, agrees
with f on sets A \in sets F.
      fix A assume A-in-sets-F: A \in sets F
     have AE \ x \ in \ M. \ norm \ (indicator \ A \ x *_R \ cond-exp \ M \ F \ (s \ (r \ n)) \ x) \leq cond-exp
M F (\lambda x. 2 * norm (f x)) x  for n
      proof -
        have AE x in M. norm (indicator A x *_{R} cond-exp M F (s (r n)) x) < norm
(cond\text{-}exp\ M\ F\ (s\ (r\ n))\ x) unfolding indicator\text{-}def by simp
         thus ?thesis using cond-exp-bounded[of n] by force
      hence lim-cond-exp-int: (\lambda n. \ LINT \ x:A|M. \ cond-exp \ M \ F \ (s \ (r \ n)) \ x)
LINT x:A|M. lim(\lambda n. cond-exp M F(s(r n)) x)
       using ae-lim-cond-exp measurable-from-subalg [OF subalg borel-measurable-indicator,
OF A-in-sets-F | cond-exp-bounded
          unfolding set-lebesque-integral-def
       \textbf{by } (intro\ integral-dominated-convergence [OF\ borel-measurable-scaleR\ borel-measurable-s
integrable-cond-exp]) (fastforce simp add: tendsto-scaleR)+
      have AE x in M. norm (indicator A x *_R s (r n) x) \le 2 * norm (f x) for n
      proof -
           have AE x in M. norm (indicator A x *_R s (r n) x) \leq norm (s (r n) x)
unfolding indicator-def by simp
         thus ?thesis using assms(5)[of - r n] by fastforce
      hence lim-s-int: (\lambda n. \ LINT \ x:A|M. \ s \ (r \ n) \ x) \longrightarrow LINT \ x:A|M. \ f \ x
       using measurable-from-subalq[OF subalq borel-measurable-indicator, OF A-in-sets-F]
LIMSEQ-subseq-LIMSEQ[OF\ assms(4)\ strict-mono-r]\ assms(5)
```

have LINT x:A|M.  $lim\ (\lambda n.\ cond\text{-}exp\ M\ F\ (s\ (r\ n))\ x) = lim\ (\lambda n.\ LINT\ x:A|M.\ cond\text{-}exp\ M\ F\ (s\ (r\ n))\ x)$  using  $limI[OF\ lim\text{-}cond\text{-}exp\text{-}int]$  by argo

also have ... =  $lim (\lambda n. LINT x: A|M. s(r n) x)$  using has-cond-expD(1)[OF

 $by \ (intro\ integral-dominated-convergence [OF\ borel-measurable-scale R\ borel-measurable-s$ 

unfolding set-lebesgue-integral-def comp-def

integrable-2f]) (fastforce simp add: tendsto-scaleR)+

```
has\text{-}cond\text{-}exp\text{-}simple[OF\ assms(2,3)]\ A\text{-}in\text{-}sets\text{-}F,\ symmetric}] by presburger
   also have ... = LINT x:A|M. fx using limI[OF lim-s-int] by argo
   finally have LINT x:A|M. lim(\lambda n. cond\text{-}exp\ M\ F\ (s\ (r\ n))\ x) = LINT\ x:A|M.
fx.
 }
Putting it all together, we have the statement we are looking for.
  hence has-cond-exp M F f (\lambda x. lim (\lambda i. cond-exp M F (s (r i)) x)) using
assms(1) lim-integrable by (intro has-cond-expI', auto)
 thus thesis using AE-Cauchy Cauchy-convergent strict-mono-r by (auto intro!:
qed
Now, we can show that the conditional expectation is well-defined for all
integrable functions.
corollary has-cond-expI:
 fixes f :: 'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology,banach}\}
 assumes integrable M f
 shows has-cond-exp M F f (cond-exp M F f)
proof -
 obtain s where s-is: \bigwedge i. simple-function M (s i) \bigwedge i. emeasure M {y \in space\ M.
\{s \mid y \neq 0\} \neq \infty \land x. \ x \in space \ M \Longrightarrow (\lambda i. \ s \mid x) \longrightarrow fx \land x \ i. \ x \in space \ M \Longrightarrow \{\lambda i. \ s \mid x\} 
norm\ (s\ i\ x) \le 2*norm\ (f\ x) using integrable-implies-simple-function-sequence [OF]
assms] by blast
 show ?thesis using has-cond-exp-simple-lim[OF assms s-is] has-cond-exp-charact(1)
by metis
qed
4.2
       Properties
The defining property of the conditional expectation now always holds, given
that the function f is integrable.
lemma cond-exp-set-integral:
 fixes f :: 'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology,banach\}
 assumes integrable M f A \in sets F
 shows (\int x \in A. fx \partial M) = (\int x \in A. cond\text{-}exp M F fx \partial M)
 using has-cond-expD(1)[OF\ has-cond-expI,\ OF\ assms] by argo
The following property of the conditional expectation is called the "Tower
Property".
lemma cond-exp-nested-subalg:
```

has-cond-exp-charact(2),  $THEN\ AE\text{-}symmetric]\ sigma\text{-}finite\text{-}subalgebra.has\text{-}cond\text{-}expI[OF]$ 

shows  $AE \xi$  in M. cond-exp  $M F f \xi = cond$ -exp M F (cond-exp  $M G f) \xi$ 

 $sigma-finite-subalgebra.intro[OF\ assms(2)]]\ nested-subalg-is-sigma-finite)$ 

using has-cond-expI assms sigma-finite-subalgebra-def by (auto intro!: has-cond-exp-nested-subalg[THEN]

fixes  $f:: 'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology,banach}\}$  assumes integrable M f subalgebra M G subalgebra G F

The conditional expectation is linear.

```
lemma cond-exp-add:
  fixes f :: 'a \Rightarrow 'b::\{second\text{-}countable\text{-}topology,banach}\}
  assumes integrable M f integrable M g
  shows AE \ x in M. cond\text{-}exp \ M \ F \ (\lambda x. \ f \ x + g \ x) \ x = cond\text{-}exp \ M \ F \ f \ x + g \ x
cond-exp M F g x
 using has-cond-exp-add OF has-cond-expI(1,1), OF assms, THEN has-cond-exp-charact (2)
lemma cond-exp-diff:
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
  assumes integrable M f integrable M q
  shows AE x in M. cond-exp M F (\lambda x. f x - g x) x = cond-exp M F f x -
cond-exp M F g x
 using has\text{-}cond\text{-}exp\text{-}add[OF\text{-}has\text{-}cond\text{-}exp\text{-}scaleR\text{-}right, OF }has\text{-}cond\text{-}exp}I(1,1),
OF assms, THEN has-cond-exp-charact(2), of -1] by simp
lemma cond-exp-diff':
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
  assumes integrable M f integrable M g
 shows AE \ x \ in \ M. \ cond-exp \ M \ F \ (f - g) \ x = cond-exp \ M \ F \ f \ x - cond-exp \ M
  unfolding fun-diff-def using assms by (rule cond-exp-diff)
lemma cond-exp-scaleR-left:
  fixes f :: 'a \Rightarrow real
  assumes integrable M f
  shows AE \ x \ in \ M. \ cond\text{-}exp \ M \ F \ (\lambda x. \ f \ x *_R \ c) \ x = \ cond\text{-}exp \ M \ F \ f \ x *_R \ c
  {\bf using} \ cond\text{-}exp\text{-}set\text{-}integral[OF \ assms] \ subalg \ assms \ {\bf unfolding} \ subalgebra\text{-}def
  by (intro cond-exp-charact,
      subst set-integral-scaleR-left, blast, intro assms,
      subst set-integral-scaleR-left, blast, intro integrable-cond-exp)
      auto
```

The conditional expectation operator is a contraction, i.e. a bounded linear operator with operator norm less than or equal to 1.

To show this we first obtain a subsequence  $\lambda x$  *i.* s (r i) x, such that  $\lambda i$ . cond-exp M F (s (r i)) x converges to cond-exp M F f x a.e. Afterwards, we obtain a sub-subsequence  $\lambda x$  i. s (r (r' i)) x, such that  $\lambda i$ . cond-exp M F  $(\lambda x. norm$  (s (r i))) x converges to cond-exp M F  $(\lambda x. norm$  (f x)) x a.e. Finally, we show that the inequality holds by showing that the terms of the subsequences obey the inequality and the fact that a subsequence of a convergent sequence converges to the same limit.

```
lemma cond-exp-contraction:

fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach} \}

assumes integrable M f

shows AE \ x \ in \ M. \ norm \ (cond\text{-}exp \ M \ F \ f \ x) \leq cond\text{-}exp \ M \ F \ (\lambda x. \ norm \ (f \ x))
```

x proof -

**obtain** s where s:  $\bigwedge i$ . simple-function M (s i)  $\bigwedge i$ . emeasure M { $y \in space M$ . s i  $y \neq 0$ }  $\neq \infty \bigwedge x$ .  $x \in space M \Longrightarrow (\lambda i. s i x) \longrightarrow f x \bigwedge i x$ .  $x \in space M \Longrightarrow norm (s i x) \leq 2 * norm (f x)$ 

by (blast intro: integrable-implies-simple-function-sequence[OF assms])

obtain r where r: strict-mono r and has-cond-exp M F f  $(\lambda x. lim$   $(\lambda i. cond$ -exp M F (s(ri)) x)) AE x in M.  $(\lambda i. cond$ -exp M F (s(ri)) x)  $\longrightarrow lim$   $(\lambda i. cond$ -exp M F (s(ri)) x)

using has-cond-exp-simple-lim[OF assms s] unfolding convergent-LIMSEQ-iff by blast

**hence** r-tendsto: AE x in M.  $(\lambda i. cond\text{-}exp\ M\ F\ (s\ (r\ i))\ x) \longrightarrow cond\text{-}exp\ M$  F f x using has-cond-exp-charact(2) by force

have norm-s-r:  $\bigwedge i$ . simple-function M ( $\lambda x$ . norm (s (r i) x))  $\bigwedge i$ . emeasure M { $y \in space\ M$ . norm (s (r i) y)  $\neq 0$ }  $\neq \infty \bigwedge x$ .  $x \in space\ M \Longrightarrow (\lambda i$ . norm (s (r i) x))  $\longrightarrow$  norm (f x)  $\bigwedge i$  x.  $x \in space\ M \Longrightarrow$  norm (norm (s (r i) x))  $\leq 2 *$  norm (norm (f x))

using s by (auto intro: LIMSEQ-subseq-LIMSEQ[OF tendsto-norm r, unfolded comp-def] simple-function-compose1)

obtain r' where r': strict-mono r' and has-cond-exp M F  $(\lambda x. norm (f x)) (\lambda x. lim (\lambda i. cond-exp <math>M$  F  $(\lambda x. norm (s (r (r' i)) x)) x)) <math>AE$  x in M.  $(\lambda i. cond-exp <math>M$  F  $(\lambda x. norm (s (r (r' i)) x)) x) <math>\longrightarrow lim (\lambda i. cond-exp M F (\lambda x. norm (s (r (r' i)) x)) x)$  using has-cond-exp-simple-lim[OF integrable-norm norm-s-r, OF assms] unfolding convergent-LIMSEQ-iff by blast

**hence** r'-tendsto:  $AE \ x \ in \ M. \ (\lambda i. \ cond\text{-}exp \ M \ F \ (\lambda x. \ norm \ (s \ (r \ (r' \ i)) \ x)) \ x)$   $\longrightarrow cond\text{-}exp \ M \ F \ (\lambda x. \ norm \ (f \ x)) \ x \ using \ has-cond\text{-}exp-charact(2) \ by \ force$ 

have  $AE \ x \ in \ M. \ \forall i. \ norm \ (cond-exp \ MF \ (s \ (r \ (r' \ i))) \ x) \le cond-exp \ MF \ (\lambda x. \ norm \ (s \ (r \ (r' \ i)) \ x)) \ x \ using \ s \ by \ (auto \ intro: \ cond-exp-contraction-simple \ simp \ add: \ AE-all-countable)$ 

ultimately show ?thesis using LIMSEQ-le r'-tendsto by fast

The following lemmas are called "pulling out whats known". We first show the statement for real-valued functions using the lemma real-cond-exp-intg, which is already present. We then show it for arbitrary g using the lecture notes of Gordan Zitkovic for the course "Theory of Probability I".

**lemma** cond-exp-measurable-mult:

```
fixes f g :: 'a \Rightarrow real
```

**assumes** [measurable]: integrable M ( $\lambda x. fx * gx$ ) integrable M  $gf \in borel$ -measurable F

```
shows integrable M (\lambda x. f x * cond\text{-}exp M F g x)

AE x in M. cond\text{-}exp M F (\lambda x. f x * g x) x = f x * cond\text{-}exp M F g x
```

```
proof-
 show integrable: integrable M (\lambda x. fx * cond\text{-}exp \ MFgx) using cond\text{-}exp\text{-}real[OF]
assms(2)] by (intro integrable-cong-AE-imp[OF real-cond-exp-intg(1), OF assms(1,3)
assms(2)[THEN\ borel-measurable-integrable]]\ measurable-from-subalq[OF\ subalq])
 interpret sigma-finite-measure restr-to-subalg M F by (rule sigma-fin-subalg)
   fix A assume asm: A \in sets F
   hence asm': A \in sets \ M  using subalg by (fastforce \ simp \ add: \ subalgebra-def)
  have set-lebesgue-integral M A (cond-exp M F (\lambda x. fx * gx)) = set-lebesgue-integral
M \ A \ (\lambda x. \ f \ x * g \ x) \ \mathbf{by} \ (simp \ add: \ cond-exp-set-integral[OF \ assms(1) \ asm])
     also have ... = set-lebesgue-integral M A (\lambda x. f x * real-cond-exp M F g
x) using borel-measurable-times [OF borel-measurable-indicator [OF asm] assms(3)]
borel-measurable-integrable [OF assms(2)] integrable-mult-indicator [OF asm' assms(1)]
by (fastforce simp add: set-lebesque-integral-def mult. assoc[symmetric] intro: real-cond-exp-intq(2)[symmetric])
    also have ... = set-lebesque-integral M A (\lambda x. f x * cond-exp M F q x) using
cond-exp-real[OF\ assms(2)]\ asm'\ borel-measurable-cond-exp'\ borel-measurable-cond-exp2
measurable-from-subalg [OF subalg assms(3)] by (auto simp add: set-lebesgue-integral-def
intro: integral-cong-AE)
   finally have set-lebesque-integral M A (cond-exp M F (\lambda x. f x * g x)) = \int x \in A.
(f x * cond\text{-}exp \ M \ F \ g \ x)\partial M.
  hence AE x in restr-to-subalg M F. cond-exp M F (\lambda x. f x * g x) x = f
x * cond\text{-}exp \ M \ F \ g \ x \ \mathbf{by} \ (intro\ density\text{-}unique\text{-}banach\ integrable\text{-}cond\text{-}exp\ integrable}
grable integrable-in-subalg subalg, measurable, simp add: set-lebesgue-integral-def
integral-subalgebra2[OF subalg] sets-restr-to-subalg[OF subalg])
  thus AE x in M. cond-exp M F (\lambda x. f x * g x) x = f x * cond-exp M F g x by
(rule\ AE-restr-to-subalg[OF\ subalg])
qed
lemma cond-exp-measurable-scaleR:
 fixes f :: 'a \Rightarrow real and g :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}
 assumes [measurable]: integrable M (\lambda x. fx *_R gx) integrable M gf \in borel-measurable
 shows integrable M (\lambda x. f x *_R cond\text{-}exp M F g x)
       AE \ x \ in \ M. \ cond-exp \ M \ F \ (\lambda x. \ f \ x *_R \ g \ x) \ x = f \ x *_R \ cond-exp \ M \ F \ g \ x
proof -
 let ?F = restr-to-subalq M F
 have subalq': subalgebra M (restr-to-subalq M F) by (metis sets-eq-imp-space-eq
sets-restr-to-subalg subalg subalgebra-def)
  fix z assume asm[measurable]: integrable M(\lambda x. zx *_R gx) z \in borel-measurable
?F
    hence asm'[measurable]: z \in borel-measurable F using measurable-in-subalg'
subalg by blast
    have integrable M (\lambda x. z x *_R cond\text{-}exp M F g x) LINT x|M. z x *_R g x =
LINT \ x | M. \ z \ x *_R \ cond-exp \ M \ F \ g \ x
   proof -
     obtain s where s-is: \bigwedge i. simple-function ?F (s i) \bigwedge x. x \in space ?F \Longrightarrow (\lambda i.
```

```
s\ i\ x) \longrightarrow z\ x \land i\ x.\ x \in space\ ?F \Longrightarrow norm\ (s\ i\ x) \le 2*norm\ (z\ x) using borel-measurable-implies-sequence-metric[OF\ asm(2)\, of\ 0]\ by\ force
```

We need to apply the dominated convergence theorem twice, therefore we need to show the following prerequisites.

have s-scaleR-g-tendsto: AE x in M. ( $\lambda i.\ s\ i\ x*_R\ g\ x) \longrightarrow z\ x*_R\ g\ x$  using s-is(2) by (simp add: space-restr-to-subalg tendsto-scaleR)

have s-scaleR-cond-exp-g-tendsto: AE x in ?F. ( $\lambda i.\ s\ i.\ x*_R\ cond-exp\ M\ F\ g$  x)  $\longrightarrow$  z x \*\_R cond-exp M F g x using s-is(2) by (simp add: tendsto-scaleR)

have s-scaleR-g-meas:  $(\lambda x.\ s\ i\ x*_R\ g\ x)\in borel$ -measurable M for i using s-is(1)[THEN borel-measurable-simple-function, THEN subalg'[THEN measurable-from-subalg]] by simp

have s-scaleR-cond-exp-g-meas:  $(\lambda x.\ s\ i\ x*_R\ cond-exp\ M\ F\ g\ x) \in borel$ -measurable ?F for i using s- $is(1)[THEN\ borel$ -measurable-simple-function] measurable-in-subalg[OF subalg\ borel-measurable-cond-exp] by (fastforce\ intro:\ borel-measurable-scaleR)

have s-scaleR-g-AE-bdd: AE x in M. norm (s i  $x *_R g x$ )  $\leq 2 * norm$  (z  $x *_R g x$ ) for i using s-is(3) by (fastforce simp add: space-restr-to-subalg mult.assoc[symmetric] mult-right-mono)

{
 fix i

have asm: integrable M ( $\lambda x$ . norm (z x) \* norm (g x)) using asm(1)[THEN integrable-norm] by simp

have  $AE \ x \ in \ ?F. \ norm \ (s \ i \ x *_R \ cond-exp \ M \ F \ g \ x) \le 2 * norm \ (z \ x) * norm \ (cond-exp \ M \ F \ g \ x) \ using \ s-is(3) \ by \ (fastforce \ simp \ add: \ mult-mono)$ 

moreover have  $AE \ x \ in \ ?F. \ norm \ (z \ x) * cond-exp \ MF \ (\lambda x. \ norm \ (g \ x)) \ x = cond-exp \ MF \ (\lambda x. \ norm \ (z \ x) * norm \ (g \ x)) \ x \ \mathbf{by} \ (rule \ cond-exp-measurable-mult(2)[THEN \ AE-symmetric, \ OF \ asm \ integrable-norm, \ OF \ assms(2), \ THEN \ AE-restr-to-subalg2[OF \ subalg]], \ auto)$ 

ultimately have  $AE \ x \ in \ ?F. \ norm \ (s \ i \ x *_R \ cond-exp \ M \ F \ g \ x) \le 2 * cond-exp \ M \ F \ (\lambda x. \ norm \ (z \ x *_R \ g \ x)) \ x \ using \ cond-exp-contraction[OF \ assms(2), THEN \ AE-restr-to-subalg2[OF \ subalg]] \ order-trans[OF - mult-mono] \ by \ fastforce$ 

**note** s-scaleR-cond-exp-g-AE-bdd = this

In the following section we need to pay attention to which measures we are using for integration. The rhs is F-measurable while the lhs is only M-measurable.

 $\{$  fix i

have s-meas-M[measurable]:  $s \ i \in borel$ -measurable M by (meson borel-measurable-simple-function measurable-from-subalg s-is(1) subalg')

have s-meas-F[measurable]:  $s \ i \in borel$ -measurable F by  $(meson\ borel$ -measurable-simple-function measurable-in-subalg' s-is $(1)\ subalg)$ 

**have** s-scaleR-eq: s i  $x *_R h$   $x = (\sum y \in s$  i 'space M. (indicator (s i - ' $\{y\} \cap space M$ )  $x *_R y$ )  $*_R h$  x) if  $x \in space M$  for x and  $h :: 'a \Rightarrow 'b$ 

using simple-function-indicator-representation[OF s-is(1), of x i] that unfolding  $space-restr-to-subalg\ scaleR-left.sum[of - - h\ x,\ symmetric]$  by presburger

```
have LINT x|M. s i x *_R g x = LINT x|M. (\sum y \in s i 'space M. in-
dicator (s \ i - `\{y\} \cap space \ M) \ x *_R y *_R g x) using s-scaleR-eq by (intro
Bochner-Integration.integral-cong) auto
           also have ... = (\sum y \in s \ i \ 'space \ M. \ LINT \ x|M. \ indicator \ (s \ i \ -'
\{y\} \cap space M) \ x *_R y *_R g x) \ \mathbf{by} \ (intro \ Bochner-Integration.integral-sum \ in-
tegrable-mult-indicator[OF-integrable-scaleR-right]\ assms(2))\ simp
       also have ... = (\sum y \in s \ i \ 'space \ M. \ y *_R set-lebesgue-integral \ M \ (s \ i - '
\{y\} \cap space\ M)\ g)\ \mathbf{by}\ (simp\ only:\ set-lebesgue-integral-def[symmetric])\ simp
      also have ... = (\sum y \in s \ i \ 'space \ M. \ y *_R set-lebesgue-integral \ M \ (s \ i - ' \{y\}
\cap space M) (cond-exp M F g)) using assms(2) subalg borel-measurable-vimage[OF]
s-meas-F] by (subst cond-exp-set-integral, auto simp add: subalgebra-def)
      also have ... = (\sum y \in s \ i \ `space M. \ LINT \ x | M. \ indicator \ (s \ i - `\{y\} \cap space \ )
M) x *_R y *_R cond-exp M F g x by (simp only: set-lebesgue-integral-def[symmetric])
simp
      also have ... = LINT x|M. (\sum y \in s \ i 'space M. indicator (s \ i - '\{y\} \cap space
M) x *_R y *_R cond-exp M F g x by (intro Bochner-Integration.integral-sum[symmetric]
integrable-mult-indicator[OF - integrable-scaleR-right]) auto
       also have ... = LINT x|M. s i x *_R cond-exp M F g x using s-scaleR-eq
by (intro Bochner-Integration.integral-cong) auto
       finally have LINT x|M. s i x *_R g x = LINT x|?F. s i x *_R cond-exp M F
g \ x \ \mathbf{by} \ (simp \ add: integral-subalgebra2[OF \ subalg])
```

Now we just plug in the results we obtained into DCT, and use the fact that limits are unique.

**note** integral-s-eq = this

show integrable M ( $\lambda x. zx *_R cond-exp\ M\ F\ g\ x$ ) using s-scaleR-cond-exp-g-meas asm(2) borel-measurable-cond-exp' by (intro integrable-from-subalg[OF subalg] integrable-cond-exp integrable-dominated-convergence[OF - - - s-scaleR-cond-exp-g-tendsto s-scaleR-cond-exp-g-AE-bdd]) (auto intro: measurable-from-subalg[OF subalg] integrable-in-subalg measurable-in-subalg subalg)

```
note * = this
The main statement now follows with z = (\lambda x. indicat-real A x * f x).
  show integrable M (\lambda x. f x *_R cond\text{-}exp M F g x) using * assms measur-
able-in-subalg[OF\ subalg] by blast
   fix A assume asm: A \in F
    hence integrable M (\lambda x. indicat-real A x *_R f x *_R g x) using subalg by
(fastforce simp add: subalgebra-def intro!: integrable-mult-indicator assms(1))
    hence set-lebesgue-integral M A (\lambda x. f x *_R g x) = set-lebesgue-integral M A
(\lambda x. f x *_R cond\text{-}exp M F g x) unfolding set-lebesgue-integral-def using asm by
(auto intro!: * measurable-in-subalg[OF subalg])
 thus AE \ x \ in \ M. \ cond-exp \ M \ F \ (\lambda x. \ f \ x *_R g \ x) \ x = f \ x *_R \ cond-exp \ M \ F \ g \ x
using borel-measurable-cond-exp by (intro cond-exp-charact, auto intro!: * assms
measurable-in-subalg[OF\ subalg])
qed
lemma cond-exp-sum [intro, simp]:
  fixes f :: 't \Rightarrow 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology,banach\}
 assumes [measurable]: \land i. integrable M (f i)
 shows AE \ x \ in \ M. \ cond\text{-}exp \ M \ F \ (\lambda x. \sum i \in I. \ f \ i \ x) \ x = (\sum i \in I. \ cond\text{-}exp \ M \ F
(f i) x)
proof (rule has-cond-exp-charact, intro has-cond-expI')
  fix A assume [measurable]: A \in sets F
 then have A-meas [measurable]: A \in sets M by (meson subsetD subalg subalge-
bra-def)
  have (\int x \in A. (\sum i \in I. \ f \ i \ x) \partial M) = (\int x. (\sum i \in I. \ indicator \ A \ x *_R f \ i \ x) \partial M)
unfolding set-lebesgue-integral-def by (simp add: scaleR-sum-right)
 also have ... = (\sum i \in I. (\int x. indicator A x *_R f i x \partial M)) using assms by (auto
intro!: Bochner-Integration.integral-sum\ integrable-mult-indicator)
 also have ... = (\sum i \in I. (\int x. indicator A x *_R cond-exp M F (f i) x \partial M)) using
cond-exp-set-integral[OF assms] by (simp add: set-lebesgue-integral-def)
  also have ... = (\int x. (\sum i \in I. indicator \ A \ x *_R cond-exp \ M \ F \ (f \ i) \ x)\partial M)
using assms by (auto intro!: Bochner-Integration.integral-sum[symmetric] inte-
grable-mult-indicator)
 also have ... = (\int x \in A. (\sum i \in I. cond\text{-}exp \ M \ F \ (f \ i) \ x) \partial M) unfolding set-lebesgue-integral-def
by (simp add: scaleR-sum-right)
 finally show (\int x \in A. (\sum i \in I. fi x) \partial M) = (\int x \in A. (\sum i \in I. cond-exp M F (fi))
(x)\partial M) by auto
qed (auto simp add: assms integrable-cond-exp)
```

## 4.3 Linearly Ordered Banach Spaces

In this subsection we show monotonicity results concerning the conditional expectation operator.

**lemma** cond-exp-qr-c:

```
fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or\text{-}
dered-real-vector}
 assumes integrable M f AE x in M. f x > c
 shows AE x in M. cond-exp M F f x > c
proof -
  define X where X = \{x \in space M. cond-exp M F f x \leq c\}
  have [measurable]: X \in sets \ F \ unfolding \ X-def \ by \ measurable \ (metis \ sets.top)
subalg\ subalgebra-def)
  hence X-in-M: X \in sets \ M using sets-restr-to-subalg subalgebra-def by
blast
 have emeasure M X = 0
 proof (rule ccontr)
   assume emeasure M X \neq 0
   have emeasure (restr-to-subalg M F) X = emeasure M X by (simp add: emea-
sure-restr-to-subalq subalq)
   hence emeasure (restr-to-subala M F) X > 0 using \langle \neg (emeasure M X) = 0 \rangle
gr-zeroI by auto
    then obtain A where A: A \in sets (restr-to-subalg M F) A \subseteq X emeasure
(restr-to-subalg M F) A > 0 emeasure (restr-to-subalg M F) A < \infty
   using sigma-fin-subalg by (metis emeasure-notin-sets enrical-0 infinity-enrical-def
le-less-linear\ neq-top-trans\ not-gr-zero\ order-refl\ sigma-finite-measure.approx-PInf-emeasure-with-finite)
   hence [simp]: A \in sets \ F using subalg sets-restr-to-subalg by blast
   hence A-in-sets-M[simp]: A \in sets \ M using sets-restr-to-subalg subalg subalg
gebra-def by blast
    have [simp]: set-integrable M A (\lambda x. c) using A subalg by (auto simp add:
set-integrable-def emeasure-restr-to-subalg)
   have [simp]: set-integrable M A f unfolding set-integrable-def by (rule inte-
grable-mult-indicator, auto\ simp\ add:\ assms(1))
   have AE \ x \ in \ M. \ indicator \ A \ x *_R \ c = indicator \ A \ x *_R \ f \ x
   proof (rule integral-eq-mono-AE-eq-AE)
   have (\int x \in A. \ c \ \partial M) \le (\int x \in A. \ f \ x \ \partial M) using assms(2) by (intro\ set\ -integral\ -mono\ -AE\ -banach)
auto
     moreover
          have (\int x \in A. \ f \ x \ \partial M) = (\int x \in A. \ cond\text{-}exp \ M \ F \ f \ x \ \partial M) by (rule
cond-exp-set-integral, auto simp add: assms)
     also have ... \leq (\int x \in A. \ c \ \partial M) using A by (auto intro!: set-integral-mono-banach
simp add: X-def)
       finally have (\int x \in A. \ f \ x \ \partial M) \le (\int x \in A. \ c \ \partial M) by simp
     ultimately show LINT x|M. indicator A \times_R c = LINT \times_R M. indicator A
x *_R f x unfolding set-lebesgue-integral-def by simp
    show AE x in M. indicator A x *_R c \leq indicator A x *_R f x using assms by
(auto simp add: X-def indicator-def)
   qed (auto simp add: set-integrable-def[symmetric])
   hence AE \ x \in A \ in \ M. \ c = f \ x \ by \ auto
   hence AE \ x \in A \ in \ M. \ False \ using \ assms(2) by auto
   hence A \in null-sets M using AE-iff-null-sets A-in-sets-M by metis
    thus False using A(3) by (simp add: emeasure-restr-to-subalg null-setsD1
```

```
subalg
 qed
 thus ?thesis using AE-iff-null-sets[OF X-in-M] unfolding X-def by auto
corollary cond-exp-less-c:
  fixes f::'a \Rightarrow 'b:: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes integrable M f AE x in M. f x < c
 shows AE x in M. cond-exp M F f x < c
proof -
 have AE x in M. cond-exp M F f x = - cond-exp M F (\lambda x. - f x) x using
cond-exp-uminus[OF assms(1)] by auto
 moreover have AE x in M. cond-exp M F (\lambda x. - f x) x > - c using assms
by (intro cond-exp-qr-c) auto
 ultimately show ?thesis by (force simp add: minus-less-iff)
qed
lemma cond-exp-mono-strict:
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes integrable M f integrable M g AE x in M. f x < g x
 shows AE x in M. cond-exp M F f x < cond-exp M F g x
 using cond-exp-less-c[OF\ Bochner-Integration.integrable-diff, OF\ assms(1,2), of
\theta
       cond-exp-diff[OF assms(1,2)] assms(3) by auto
lemma cond-exp-ge-c:
  fixes f::'a \Rightarrow 'b:: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes [measurable]: integrable M f
    and AE x in M. f x > c
 shows AE x in M. cond-exp M F f x \ge c
proof -
 let ?F = restr-to-subalg M F
 interpret sigma-finite-measure restr-to-subala M F using sigma-fin-subala by
auto
   fix A assume asm: A \in sets ?F 0 < measure ?F A
  have [simp]: sets ?F = sets\ F\ measure\ ?F\ A = measure\ M\ A\ using\ asm\ by\ (auto
simp\ add: measure-def\ sets-restr-to-subalg[OF\ subalg]\ emeasure-restr-to-subalg[OF\ subalg]
subalg()
   have M-A: emeasure M A < \infty using measure-zero-top asm by (force simp
add: top.not-eq-extremum)
   hence F-A: emeasure ?F A < \infty using asm(1) emeasure-restr-to-subalg subalg
by fastforce
```

have set-lebesgue-integral M A  $(\lambda$ -.  $c) \leq$  set-lebesgue-integral M A f using assms asm M-A subalg by (intro set-integral-mono-AE-banach, auto simp add: set-integrable-def integrable-mult-indicator subalgebra-def sets-restr-to-subalg)

```
also have ... = set-lebesque-integral M A (cond-exp M F f) using cond-exp-set-integral OF
assms(1)] asm by auto
  also have \dots = set-lebesgue-integral ?F A (cond-exp M F f) unfolding set-lebesgue-integral-def
using asm borel-measurable-cond-exp by (intro integral-subalgebra2 OF subalg, sym-
metric, simp)
  finally have (1 \mid measure ?FA) *_R set-lebesgue-integral ?FA (cond-exp M F f)
\in \{c..\} using asm subalq M-A by (auto simp add: set-integral-const subalgebra-def
intro!: pos-divideR-le-eq[THEN iffD1])
 }
  thus ?thesis using AE-restr-to-subalg[OF subalg] averaging-theorem[OF inte-
grable-in-subalg closed-atLeast, OF subalg borel-measurable-cond-exp integrable-cond-exp
by auto
qed
corollary cond-exp-le-c:
  fixes f: 'a \Rightarrow 'b: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes integrable M f
     and AE x in M. f x \leq c
 shows AE x in M. cond-exp M F f x \le c
proof -
  have AE x in M. cond-exp M F f x = - cond-exp M F (\lambda x. - f x) x using
cond-exp-uminus[OF assms(1)] by force
  moreover have AE x in M. cond-exp M F (\lambda x. - f x) x \ge -c using assms
by (intro cond-exp-ge-c) auto
 ultimately show ?thesis by (force simp add: minus-le-iff)
qed
corollary cond-exp-mono:
  fixes f::'a \Rightarrow 'b:: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes integrable M f integrable M g AE x in M. f x \leq g x
 shows AE x in M. cond-exp M F f x \leq cond-exp M F g x
  using cond-exp-le-c[OF\ Bochner-Integration.integrable-diff, OF\ assms(1,2), of
\theta
       cond-exp-diff[OF assms(1,2)] assms(3) by auto
corollary cond-exp-min:
  fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes integrable M f integrable M g
 shows AE \xi in M. cond-exp M F (\lambda x. min (f x) (g x)) \xi \leq min (cond-exp M F)
f \xi) (cond-exp M F g \xi)
proof -
 have AE \xi in M. cond-exp M F (\lambda x. min (f x) (g x)) \xi \leq cond-exp <math>M F f \xi by
(intro cond-exp-mono integrable-min assms, simp)
  moreover have AE \ \xi \ in \ M. \ cond-exp \ M \ F \ (\lambda x. \ min \ (f \ x) \ (g \ x)) \ \xi \leq cond-exp
M F g \xi by (intro cond-exp-mono integrable-min assms, simp)
  ultimately show AE \xi in M. cond-exp M F (\lambda x. min (f x) (g x)) \xi \leq min
```

```
(cond\text{-}exp\ M\ F\ f\ \xi)\ (cond\text{-}exp\ M\ F\ g\ \xi)\ \mathbf{by}\ fastforce
qed
corollary cond-exp-max:
  fixes f: 'a \Rightarrow 'b: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector}
 assumes integrable M f integrable M g
 shows AE \xi in M. cond-exp M F (\lambda x. max (f x) (g x)) \xi \geq max (cond-exp M F)
f \xi) (cond-exp M F g \xi)
proof -
 have AE \xi in M. cond-exp M F (\lambda x. max (f x) (g x)) \xi \geq cond-exp M F f \xi by
(intro cond-exp-mono integrable-max assms, simp)
 moreover have AE \xi in M. cond-exp M F (\lambda x. max (f x) (g x)) \xi \geq cond-exp
M F g \xi by (intro cond-exp-mono integrable-max assms, simp)
  ultimately show AE \xi in M. cond-exp M F (\lambda x. max (f x) (g x)) \xi \geq max
(cond\text{-}exp\ M\ F\ f\ \xi)\ (cond\text{-}exp\ M\ F\ q\ \xi)\ \mathbf{by}\ fastforce
qed
corollary cond-exp-inf:
  fixes f: 'a \Rightarrow 'b: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector, lattice}
 assumes integrable M f integrable M g
  shows AE \xi in M. cond-exp M F (\lambda x. inf (f x) (g x)) \xi \leq inf (cond-exp <math>M F f
\xi) (cond-exp M F q \xi)
  unfolding inf-min using assms by (rule cond-exp-min)
corollary cond-exp-sup:
  fixes f: 'a \Rightarrow 'b: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, or-
dered-real-vector, lattice}
 assumes integrable M f integrable M g
 shows AE \notin in M. cond-exp M F (\lambda x. sup (f x) (g x)) \notin sup (cond-exp M F f
\xi) (cond-exp M F q \xi)
 unfolding sup-max using assms by (rule cond-exp-max)
end
4.4
       Probability Spaces
lemma (in prob-space) sigma-finite-subalgebra-restr-to-subalg:
 assumes subalgebra\ M\ F
 shows sigma-finite-subalgebra M F
proof (intro sigma-finite-subalgebra.intro)
 interpret\ F:\ prob\ -space\ restr-to\ -subalg\ M\ F\ using\ assms\ prob\ -space\ -restr-to\ -subalg
prob-space-axioms by blast
 show sigma-finite-measure (restr-to-subala MF) by (rule F. sigma-finite-measure-axioms)
qed (rule assms)
lemma (in prob-space) cond-exp-trivial:
```

fixes  $f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach}\}$ 

```
assumes integrable M f
   shows AE x in M. cond-exp M (sigma (space M) \{\}) f x = expectation f
proof -
  interpret sigma-finite-subalgebra\ M\ sigma\ (space\ M)\ \{\}\  by (auto intro: sigma-finite-subalgebra-restr-to-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-subalgebra-
simp add: subalgebra-def sigma-sets-empty-eq)
 show ?thesis using assms by (intro cond-exp-charact) (auto simp add: sigma-sets-empty-eq
set-lebesque-integral-def prob-space cong: Bochner-Integration.integral-cong)
qed
The following lemma shows that independent \sigma-algebras don't matter for
the conditional expectation.
lemma (in prob-space) cond-exp-indep-subalgebra:
   fixes f: 'a \Rightarrow 'b: \{second\text{-}countable\text{-}topology, banach, real\text{-}normed\text{-}field}\}
  assumes subalgebra: subalgebra M F subalgebra M G
         and independent: indep-set G (sigma (space M) (F \cup vimage-algebra (space
M) f borel)
   assumes [measurable]: integrable M f
  shows AE \ x \ in \ M. \ cond-exp \ M \ (sigma \ (space \ M) \ (F \cup G)) \ f \ x = cond-exp \ M \ F
f x
proof -
  interpret Un-sigma: sigma-finite-subalgebra M sigma (space M) (F \cup G) using
assms(1,2) by (auto intro!: sigma-finite-subalgebra-restr-to-subalg sets.sigma-sets-subset
simp add: subalgebra-def space-measure-of-conv sets-measure-of-conv)
 interpret sigma-finite-subalgebra MF using assms by (auto intro: sigma-finite-subalgebra-restr-to-subalg)
     \mathbf{fix} \ A
     assume asm: A \in sigma \ (space \ M) \ \{a \cap b \mid a \ b. \ a \in F \land b \in G\}
       have in-events: sigma-sets (space M) \{a \cap b \mid a \ b. \ a \in sets \ F \land b \in sets \}
G \subseteq events using subalgebra by (intro sets.sigma-sets-subset, auto simp add:
subalgebra-def)
     have Int-stable \{a \cap b \mid a \ b. \ a \in F \land b \in G\}
     proof -
        {
           fix af bf ag bg
           assume F: af \in F \ bf \in F \ and \ G: ag \in G \ bg \in G
           have af \cap bf \in F by (intro\ sets.Int\ F)
           moreover have ag \cap bg \in G by (intro sets.Int G)
            ultimately have \exists a \ b. \ af \cap ag \cap (bf \cap bg) = a \cap b \wedge a \in sets \ F \wedge b \in
sets G by (metis inf-assoc inf-left-commute)
        thus ?thesis by (force intro!: Int-stableI)
     qed
       moreover have \{a \cap b \mid a \ b. \ a \in F \land b \in G\} \subseteq Pow \ (space \ M) using
subalgebra by (force simp add: subalgebra-def dest: sets.sets-into-space)
     moreover have A \in sigma-sets (space M) \{a \cap b \mid a \ b. \ a \in F \land b \in G\} using
calculation asm by force
        ultimately have set-lebesgue-integral M A f = set-lebesgue-integral M A
(cond\text{-}exp\ M\ F\ f)
     proof (induction rule: sigma-sets-induct-disjoint)
```

```
case (basic\ A)
then obtain a\ b where A: A = a \cap b\ a \in F\ b \in G by blast
```

**hence**  $events[measurable]: a \in events b \in events using subalgebra by (auto simp add: subalgebra-def)$ 

 $\mathbf{have} \ [simp] : sigma-sets \ (space \ M) \ \{indicator \ b - `A \cap space \ M \ | A. \ A \in borel\} \subseteq G$ 

 $\begin{tabular}{ll} \textbf{using} \ borel-measurable-indicator} [OF\ A(3),\ THEN\ measurable-sets] \ sets.top \\ subalgebra \end{tabular}$ 

**by** (intro sets.sigma-sets-subset') (fastforce simp add: subalgebra-def)+

**have** Un-in-sigma:  $F \cup vimage$ -algebra (space M) f borel  $\subseteq sigma$  (space M)  $(F \cup vimage$ -algebra (space M) f borel) by (metis equalityE le-supI sets.space-closed sigma-le-sets space-vimage-algebra subalg subalgebra-def)

have [intro]: indep-var borel (indicator b) borel ( $\lambda\omega$ . indicator a  $\omega*_R f\omega$ ) proof -

**have** [simp]: sigma-sets (space M)  $\{(\lambda \omega. indicator \ a \ \omega *_R f \ \omega) - `A \cap space M \ | A. \ A \in borel\} \subseteq sigma (space M) (F \cup vimage-algebra (space M) f borel)$ 

**have** \*:  $(\lambda \omega$ . indicator a  $\omega *_R f \omega) \in borel$ -measurable (sigma (space M)  $(F \cup vimage$ -algebra (space M) f borel))

 $\textbf{using} \ borel-measurable-indicator [OF\ A(2),\ THEN\ measurable-sets,\ OF\ borel-open]\ subalgebra$ 

 $\mathbf{by}\ (intro\ borel-measurable\text{-}scaleR\ borel-measurableI\ Un\text{-}in\text{-}sigma[THEN\ subsetD]})$ 

 $\stackrel{'}{(}auto\;simp\;add:\;space-measure-of-conv\;subalgebra-def\;sets-vimage-algebra2)$ 

thus ?thesis using measurable-sets[OF\*] by  $(intro\ sets.sigma\text{-}sets\text{-}subset',\ auto\ simp\ add:\ space\text{-}measure\text{-}of\text{-}conv)$ 

qed

**have** indep-set (sigma-sets (space M) {indicator b - 'A  $\cap$  space M | A. A  $\in$  borel}) (sigma-sets (space M) {( $\lambda\omega$ . indicator  $a\ \omega*_R f\ \omega$ ) - 'A  $\cap$  space M | A. A  $\in$  borel})

using independent unfolding indep-set-def by  $(rule\ indep-sets-mono-sets, auto\ split:\ bool.split)$ 

thus ?thesis by (subst indep-var-eq, auto intro!: borel-measurable-scaleR) qed

have [intro]: indep-var borel (indicator b) borel ( $\lambda\omega$ . indicat-real a  $\omega*_R$  cond-exp M F f  $\omega$ )

proof -

**have** [simp]:sigma-sets  $(space\ M)$   $\{(\lambda\omega.\ indicator\ a\ \omega*_R\ cond-exp\ M\ F\ f\ \omega)$  - ' $A\cap space\ M\ |A.\ A\in borel\}\subseteq sigma\ (space\ M)\ (F\cup vimage-algebra\ (space\ M)\ f\ borel)$ 

proof -

**have** \*:  $(\lambda \omega. indicator \ a \ \omega *_R cond-exp \ M \ F \ f \ \omega) \in borel-measurable (sigma (space M) (F \cup vimage-algebra (space M) f borel))$ 

 $\mathbf{using}\ borel-measurable\text{-}indicator[OF\ A(2),\ THEN\ measurable\text{-}sets,\ OF$ 

```
borel-open] subalgebra
                                 borel-measurable-cond-exp[THEN measurable-sets, OF borel-open, of
- M F f]
                        by (intro borel-measurable-scaleR borel-measurableI Un-in-sigma[THEN
subsetD)
                            (auto simp add: space-measure-of-conv subalgebra-def)
                thus ?thesis using measurable-sets[OF *] by (intro sets.sigma-sets-subset',
auto simp add: space-measure-of-conv)
               ged
              have indep-set (sigma-sets (space M) {indicator b - A \cap space M \mid A. A \in A \cap space M \mid A. A \in A \cap space M \mid A. A \in A \cap space M \mid A. A \cap sp
borel}) (sigma-sets (space M) \{(\lambda \omega. indicator \ a \ \omega *_R cond-exp \ M \ F \ f \ \omega) - `A \cap A \cap B \}
space\ M\ |A.\ A \in borel\}
                using independent unfolding indep-set-def by (rule indep-sets-mono-sets,
auto split: bool.split)
              thus ?thesis by (subst indep-var-eq, auto intro!: borel-measurable-scaleR)
           have set-lebesgue-integral M A f = (LINT \ x|M. \ indicator \ b \ x * (indicator \ a
               unfolding set-lebesque-integral-def A indicator-inter-arith
          \mathbf{by}\ (intro\ Bochner-Integration.integral-cong,\ auto\ simp\ add:\ scaleR-scaleR[symmetric]
indicator-times-eq-if(1))
           also have ... = (LINT \ x|M. \ indicator \ b \ x) * (LINT \ x|M. \ indicator \ a \ x*_R f
x)
               by (intro indep-var-lebesgue-integral
                                   Bochner-Integration.integrable-bound[OF integrable-const[of 1 :: 'b]
borel-measurable-indicator]
                                        integrable-mult-indicator[OF - assms(4)], blast) (auto simp\ add:
indicator-def)
               also have ... = (LINT \ x|M. \ indicator \ b \ x) * (LINT \ x|M. \ indicator \ a \ x*_R
cond-exp M F f x)
             using cond-exp-set-integral [OF assms(4) A(2)] unfolding set-lebesgue-integral-def
by argo
                also have ... = (LINT x|M. indicator b \ x * (indicator \ a \ x *_R \ cond-exp \ M
Ff(x)
               by (intro indep-var-lebesque-integral[symmetric]
                                   Bochner-Integration.integrable-bound[OF integrable-const[of 1 :: 'b]
borel-measurable-indicator
                                integrable-mult-indicator[OF - integrable-cond-exp], blast) (auto simp
add: indicator-def)
           also have ... = set-lebesgue-integral M A (cond-exp M F f)
               \mathbf{unfolding} set-lebesgue-integral-def A indicator-inter-arith
          by (intro Bochner-Integration.integral-cong, auto simp add: scaleR-scaleR[symmetric]
indicator-times-eq-if(1)
           finally show ?case.
        next
           case empty
           then show ?case unfolding set-lebesque-integral-def by simp
       next
```

```
case (compl A)
   have A-in-space: A \subseteq space \ M using compl using in-events sets.sets-into-space
by blast
    have set-lebesque-integral M (space M-A) f=set-lebesque-integral M (space
M-A\cup A) f-set-lebesgue-integral M A f
      using compl(1) in-events
      by (subst set-integral-Un[of space M - A A], blast)
            (simp \mid intro integrable-mult-indicator [folded set-integrable-def, OF -
assms(4)], fast)+
    also have ... = set-lebesgue-integral M (space M - A \cup A) (cond-exp M F f)
- set-lebesgue-integral M A (cond-exp M F f)
     using cond-exp-set-integral [OF assms(4) sets.top] compl subalgebra by (simp
add: subalgebra-def Un-absorb2[OF A-in-space])
     also have ... = set-lebesgue-integral M (space M - A) (cond-exp M F f)
      using compl(1) in-events
      by (subst set-integral-Un[of space M-AA], blast)
            (simp | intro integrable-mult-indicator folded set-integrable-def, OF -
integrable-cond-exp[, fast)+
     finally show ?case.
   next
     case (union A)
    have set-lebesgue-integral M (\bigcup (range A)) f = (\sum i. set-lebesgue-integral M
(A \ i) \ f)
       using union in-events
     by (intro lebesque-integral-countable-add) (auto simp add: disjoint-family-onD
intro!: integrable-mult-indicator[folded set-integrable-def, OF - assms(4)])
     also have ... = (\sum i. set-lebesgue-integral M (A i) (cond-exp M F f)) using
union by presburger
     also have ... = set-lebesgue-integral M (\bigcup (range \ A)) (cond-exp M \ F \ f)
      using union in-events
       by (intro lebesgue-integral-countable-add[symmetric]) (auto simp add: dis-
joint-family-onD intro!: integrable-mult-indicator[folded set-integrable-def, OF - in-
tegrable-cond-exp])
    finally show ?case.
   qed
 moreover have sigma\ (space\ M)\ \{a\cap b\mid a\ b.\ a\in F\land b\in G\}=sigma\ (space\ M)
M) (F \cup G)
 proof -
   have sigma-sets (space\ M) {a \cap b \mid a\ b.\ a \in sets\ F \land b \in sets\ G} = sigma-sets
(space\ M)\ (sets\ F\cup sets\ G)
   proof -
     {
      fix a b assume asm: a \in F b \in G
      hence a \cap b \in sigma\text{-}sets (space M) (F \cup G) using subalgebra unfolding
Int-range-binary by (intro sigma-sets-Inter[OF - binary-in-sigma-sets]) (force simp
add: subalgebra-def dest: sets.sets-into-space)+
     }
     moreover
```

```
{
              \mathbf{fix} \ a
              assume a \in sets F
              hence a \in sigma\text{-}sets (space M) \{a \cap b \mid a b.\ a \in sets\ F \land b \in sets\ G\}
                  using subalgebra sets.top[of G] sets.sets-into-space[of - F]
                  by (intro sigma-sets.Basic, auto simp add: subalgebra-def)
           moreover
           {
              fix a assume a \in sets \ F \lor a \in sets \ G \ a \notin sets \ F
              hence a \in sets G by blast
              hence a \in sigma\text{-}sets (space M) \{a \cap b \mid a b. a \in sets F \land b \in sets G\}
                  using subalgebra sets.top[of F] sets.sets-into-space[of - G]
                  by (intro sigma-sets.Basic, auto simp add: subalgebra-def)
           ultimately show ?thesis by (intro sigma-sets-eqI) auto
       qed
        thus ?thesis using subalgebra by (intro sigma-eqI) (force simp add: subalge-
bra-def dest: sets.sets-into-space)+
    moreover have (cond\text{-}exp\ M\ F\ f) \in borel\text{-}measurable} (sigma\ (space\ M)\ (sets\ F\ f))
\cup sets G))
    proof -
          have F \subseteq sigma \ (space \ M) \ (F \cup G) by (metis Un-least Un-upper1 mea-
sure-of-of-measure\ sets.space-closed\ sets-measure-of\ sigma-sets-subseteq\ subalq\ sub-of-of-measure
algebra(2) subalgebra-def)
     thus ?thesis using borel-measurable-cond-exp[THEN measurable-sets, OF borel-open,
of - MFf | subalgebra by (intro borel-measurable I, force simp only: space-measure-of-conv
subalgebra-def)
   qed
  ultimately show ?thesis using assms(4) integrable-cond-exp by (intro Un-sigma.cond-exp-charact)
presburger+
qed
If a random variable is independent of a \sigma-algebra F, its conditional expec-
tation cond-exp M F f is just its expectation.
lemma (in prob-space) cond-exp-indep:
   fixes f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, real\text{-}normed\text{-}field}\}
   assumes subalgebra: subalgebra M F
           and independent: indep-set \ F \ (vimage-algebra \ (space \ M) \ f \ borel)
           and integrable: integrable M f
   shows AE x in M. cond-exp M F f x = expectation f
proof -
  have indep\text{-}set\ F\ (sigma\ (space\ M)\ (sigma\ (space\ M)\ \{\}\cup (vimage\text{-}algebra\ (space\ M)\ \{\}\cup (vimage\ M)\ \{\}\cup (vimage\ M)\ \{\}\cup (vimage\ M)\ \{\}\cup (vimage\ M)\ \{\}\cup 
M) f borel)))
       using independent unfolding indep-set-def
       by (rule indep-sets-mono-sets, simp add: bool.split)
         (\it metis\ bot. extremum\ dual-order. refl\ sets. sets-measure-of-eq\ sets. sigma-sets-subset')
sets-vimage-algebra-space space-vimage-algebra sup.absorb-iff2)
```

```
hence cond-exp-indep: AE \ x \ in \ M. \ cond-exp \ M \ (sigma \ (space \ M) \ (sigma \ (space \ M))
M) {} \cup F)) f x = expectation <math>f
    {\bf using} \ cond\text{-}exp\text{-}indep\text{-}subalgebra[OF\text{-}subalgebra\text{-}integrable, of sigma (space)]}
M) {}| cond-exp-trivial[OF integrable]
   by (auto simp add: subalgebra-def sigma-sets-empty-eq)
 have sets (sigma (space M) (sigma (space M) \{\} \cup F\}) = F
   using subalgebra\ sets.top[of\ F] unfolding subalgebra-def
   by (simp add: sigma-sets-empty-eq, subst insert-absorb[of space M F], blast)
      (metis insert-absorb[OF sets.empty-sets] sets.sets-measure-of-eq)
 hence AE \ x \ in \ M. \ cond-exp \ M \ (sigma \ (space \ M) \ (sigma \ (space \ M) \ \{\} \cup F)) \ f
x = cond\text{-}exp \ M \ F \ f \ x \ by \ (rule \ cond\text{-}exp\text{-}sets\text{-}cong)
 thus ?thesis using cond-exp-indep by force
qed
end
theory Filtered-Measure
 imports\ HOL-Probability. Conditional-Expectation
begin
5
      Filtered Measure Spaces
       Filtered Measure
locale filtered-measure =
 fixes M F and t_0 :: \{second\text{-}countable\text{-}topology, order\text{-}topology, } t2\text{-}space\}
 assumes subalgebras: \bigwedge i. t_0 \leq i \Longrightarrow subalgebra\ M\ (F\ i)
     and sets-F-mono: \bigwedge i \ j. t_0 \le i \Longrightarrow i \le j \Longrightarrow sets \ (F \ i) \le sets \ (F \ j)
begin
lemma space-F[simp]:
 assumes t_0 \leq i
 shows space (F i) = space M
 using subalgebras assms by (simp add: subalgebra-def)
lemma subalgebra-F[intro]:
 assumes t_0 \leq i \ i \leq j
 shows subalgebra (F j) (F i)
 unfolding subalgebra-def using assms by (simp add: sets-F-mono)
{f lemma}\ borel-measurable-mono:
 assumes t_0 \leq i \ i \leq j
 shows borel-measurable (F i) \subseteq borel-measurable (F j)
 unfolding subset-iff by (metis assms subalgebra-F measurable-from-subalg)
end
```

locale linearly-filtered-measure = filtered-measure  $M F t_0$  for M and F :: - ::

```
\{linorder\text{-}topology\} \Rightarrow - and t_0

\mathbf{locale}\ nat\text{-}filtered\text{-}measure} = linearly\text{-}filtered\text{-}measure}\ M\ F\ 0\ \mathbf{for}\ M\ \mathbf{and}\ F::\ nat

\Rightarrow - \mathbf{locale}\ real\text{-}filtered\text{-}measure} = linearly\text{-}filtered\text{-}measure}\ M\ F\ 0\ \mathbf{for}\ M\ \mathbf{and}\ F::\ real

\Rightarrow -
```

## 5.2 $\sigma$ -Finite Filtered Measure

The locale presented here is a generalization of the *sigma-finite-subalgebra* for a particular filtration.

### 5.3 Finite Filtered Measure

```
locale\ finite-filtered-measure\ =\ filtered-measure\ +\ finite-measure
```

```
sublocale finite-filtered-measure \subseteq sigma-finite-filtered-measure using subalgebras by (unfold-locales, blast, meson dual-order.refl finite-measure-axioms finite-measure-def finite-measure-restr-to-subalg sigma-finite-measure.sigma-finite-countable)
```

```
\begin{tabular}{l} \textbf{locale} & \textit{nat-finite-filtered-measure} & \textit{m} & \textit{F} & 0 :: \textit{nat} & \textbf{for} & \textit{M} & \textit{F} \\ \textbf{locale} & \textit{real-finite-filtered-measure} & = \textit{finite-filtered-measure} & \textit{M} & \textit{F} & 0 :: \textit{real} & \textbf{for} & \textit{M} & \textit{F} \\ \end{tabular}
```

```
sublocale nat-finite-filtered-measure \subseteq nat-sigma-finite-filtered-measure .. sublocale real-finite-filtered-measure \subseteq real-sigma-finite-filtered-measure ...
```

# 5.4 Constant Filtration

```
lemma filtered-measure-constant-filtration:
assumes subalgebra M F
shows filtered-measure M (\lambda-. F) t_0
```

```
using assms by (unfold-locales) blast+
\mathbf{sublocale}\ sigma\text{-}finite\text{-}subalgebra\subseteq constant\text{-}filtration:\ sigma\text{-}finite\text{-}filtered\text{-}measure
M \lambda- :: 't :: {second-countable-topology, linorder-topology}. F t_0
 using subalg by (unfold-locales) blast+
{\bf lemma}~({\bf in}~\textit{finite-measure})~\textit{filtered-measure-constant-filtration}:
  assumes subalgebra\ M\ F
 shows finite-filtered-measure M (\lambda-. F) t_0
 using assms by (unfold-locales) blast+
end
theory Stochastic-Process
imports Filtered-Measure Measure-Space-Supplement HOL-Probability. Independent-Family
begin
      Stochastic Processes
6
       Stochastic Process
A stochastic process is a collection of random variables, indexed by a type
locale stochastic-process =
 fixes M t_0 and X :: 'b :: \{second-countable-topology, order-topology, t2-space\} \Rightarrow
'a \Rightarrow 'c :: \{second\text{-}countable\text{-}topology, banach}\}
 assumes random\text{-}variable[measurable]: \land i.\ t_0 \leq i \Longrightarrow X\ i \in borel\text{-}measurable\ M
begin
definition left-continuous where left-continuous = (AE \ \xi \ in \ M. \ \forall \ t. \ continuous
(at\text{-left }t)\ (\lambda i.\ X\ i\ \xi))
definition right-continuous where right-continuous = (AE \xi in M. \forall t. continuous)
(at\text{-}right\ t)\ (\lambda i.\ X\ i\ \xi))
end
We specify the following locales to formalize discrete time and continuous
time processes.
locale nat-stochastic-process = stochastic-process M \ 0 :: nat \ X for M \ X
locale real-stochastic-process = stochastic-process M \ 0 :: real \ X for M \ X
lemma stochastic-process-const-fun:
 assumes f \in borel-measurable M
 shows stochastic-process M t_0 (\lambda-. f) using assms by (unfold-locales)
{\bf lemma}\ stochastic\text{-}process\text{-}const:
 shows stochastic-process M t_0 (\lambda i -. c i) by (unfold-locales) simp
```

```
context stochastic-process
begin
lemma compose-stochastic:
  assumes \bigwedge i. t_0 \leq i \Longrightarrow f i \in borel\text{-}measurable\ borel
  shows stochastic-process M t_0 (\lambda i \ \xi. (f \ i) \ (X \ i \ \xi))
  by (unfold-locales) (intro measurable-compose[OF random-variable assms])
lemma norm-stochastic: stochastic-process M t_0 (\lambda i \ \xi. norm (X i \ \xi)) by (fastforce
intro: compose-stochastic)
lemma scaleR-right-stochastic:
  assumes stochastic-process\ M\ t_0\ Y
 shows stochastic-process M t_0 (\lambda i \ \xi. (Y \ i \ \xi) *_R (X \ i \ \xi))
 using stochastic-process.random-variable [OF assms] random-variable by (unfold-locales)
simp
\mathbf{lemma}\ scaleR-right-const-fun-stochastic:
  assumes f \in borel-measurable M
  shows stochastic-process M t_0 (\lambda i \ \xi. \ f \ \xi *_R (X \ i \ \xi))
  \mathbf{by}\ (unfold\text{-}locales)\ (intro\ borel\text{-}measurable\text{-}scaleR\ assms\ random\text{-}variable)
lemma scaleR-right-const-stochastic: stochastic-process M t_0 (\lambda i \ \xi. c \ i *_R (X \ i \ \xi))
  by (unfold-locales) simp
lemma add-stochastic:
  assumes stochastic-process\ M\ t_0\ Y
 shows stochastic-process M t_0 (\lambda i \ \xi. X \ i \ \xi + Y \ i \ \xi)
 using stochastic-process.random-variable [OF assms] random-variable by (unfold-locales)
simp
{f lemma} diff-stochastic:
  assumes stochastic-process M t_0 Y
 shows stochastic-process M t_0 (\lambda i \ \xi. X \ i \ \xi - Y \ i \ \xi)
 \mathbf{using}\ stochastic-process.random-variable[OF\ assms]\ random-variable\ \mathbf{by}\ (unfold-locales)
simp
lemma uminus-stochastic: stochastic-process M t_0 (-X) using scaleR-right-const-stochastic[of
\lambda-. -1] by (simp add: fun-Compl-def)
lemma partial-sum-stochastic: stochastic-process M t_0 (\lambda n \xi. \sum i \in \{t_0..n\}. X i \xi)
by (unfold-locales) simp
lemma partial-sum'-stochastic: stochastic-process M t_0 (\lambda n \xi. \sum i \in \{t_0... < n\}). X i
\xi) by (unfold-locales) simp
end
```

In the following segment, we cover basic operations on stochastic processes.

```
lemma stochastic-process-sum:

assumes \bigwedge i. i \in I \Longrightarrow stochastic-process\ M\ t_0\ (X\ i)

shows stochastic-process M\ t_0\ (\lambda k\ \xi.\ \sum i \in I.\ X\ i\ k\ \xi) using assms[THEN\ stochastic-process.random-variable] by (unfold-locales,\ auto)
```

### 6.1.1 Natural Filtration

The natural filtration induced by a stochastic process X is the filtration generated by all events involving the process up to the time index t, i.e. F  $t = \sigma(\{X \mid s \mid s. s \leq t\})$ .

**definition** natural-filtration :: 'a measure  $\Rightarrow$  'b  $\Rightarrow$  ('b  $\Rightarrow$  'a  $\Rightarrow$  'c :: topological-space)  $\Rightarrow$  'b :: {second-countable-topology, order-topology}  $\Rightarrow$  'a measure **where** natural-filtration M to  $Y = (\lambda t. family-vimage-algebra (space <math>M$ ) {Y i | i. i  $\in$  { $t_0..t$ }} borel)

```
abbreviation nat-natural-filtration \equiv \lambda M. natural-filtration M (0 :: nat) abbreviation real-natural-filtration \equiv \lambda M. natural-filtration M (0 :: real)
```

**lemma** space-natural-filtration[simp]: space (natural-filtration M  $t_0$  X t) = space M unfolding natural-filtration-def space-family-vimage-algebra ..

```
lemma sets-natural-filtration: sets (natural-filtration M t_0 X t) = sigma-sets (space M) (\bigcup i \in \{t_0..t\}. {X i - 'A \cap space M | A. A \in borel}) unfolding natural-filtration-def sets-family-vimage-algebra by (intro sigma-sets-eqI) blast+
```

```
lemma sets-natural-filtration':
 assumes borel = sigma\ UNIV\ S
 shows sets (natural-filtration M t_0 X t) = sigma-sets (space M) (\bigcup i \in \{t_0..t\}. \{X\}
i - A \cap space M \mid A. A \in S\}
proof (subst sets-natural-filtration, intro sigma-sets-eqI, clarify)
 fix i and A :: 'a set assume asm: i \in \{t_0..t\} A \in sets borel
 hence A \in sigma-sets \ UNIV \ S unfolding assms by simp
  thus X i - A \cap space M \in sigma-sets (space M) (\bigcup i \in \{t_0..t\}. \{X i - A \cap A \cap A \cap A\})
space M \mid A. A \in S \}
 proof (induction)
   case (Compl\ a)
   have Xi - (UNIV - a) \cap space M = space M - (Xi - a \cap space M) by
   then show ?case using Compl(2)[THEN sigma-sets.Compl] by presburger
  next
   case (Union \ a)
   have X \ i - ` \bigcup (range \ a) \cap space \ M = \bigcup (range \ (\lambda j. \ X \ i - ` a \ j \cap space \ M))
by blast
   then show ?case using Union(2)[THEN sigma-sets.Union] by presburger
 qed (auto intro: asm sigma-sets.Empty)
qed (intro sigma-sets.Basic, force simp add: assms)
```

**lemma** sets-natural-filtration-open:

```
sets (natural-filtration M t_0 X t) = sigma-sets (space M) (\bigcup i \in \{t_0..t\}. {X i - '
A \cap space M \mid A. open A\})
 using sets-natural-filtration' by (force simp only: borel-def mem-Collect-eq)
lemma sets-natural-filtration-oi:
 sets (natural-filtration M t_0 X t) = sigma-sets (space M) (\bigcup i \in \{t_0..t\}. {X i - 'A
\cap space M \mid A :: - :: \{linorder-topology, second-countable-topology\} set. <math>A \in range
 by (rule sets-natural-filtration'[OF borel-Ioi])
lemma sets-natural-filtration-io:
 \cap space M \mid A :: - :: \{linorder-topology, second-countable-topology\} set. <math>A \in range
lessThan\})
 by (rule sets-natural-filtration'[OF borel-Iio])
lemma sets-natural-filtration-ci:
  sets (natural-filtration M t_0 X t) = sigma-sets (space M) (\bigcup i \in \{t_0..t\}. {X i - '
A \cap space M \mid A :: real set. A \in range atLeast\}
 by (rule sets-natural-filtration'[OF borel-Ici])
context stochastic-process
begin
{f lemma}\ subalgebra-natural-filtration:
  shows subalgebra M (natural-filtration M t_0 X i)
  unfolding subalgebra-def using measurable-family-iff-sets by (force simp add:
natural-filtration-def)
lemma filtered-measure-natural-filtration:
  shows filtered-measure M (natural-filtration M t_0 X) t_0
  by (unfold-locales) (intro subalgebra-natural-filtration, simp only: sets-natural-filtration,
intro sigma-sets-subseteq, force)
In order to show that the natural filtration constitutes a filtered \sigma-finite
measure, we need to provide a countable exhausting set in the preimage of
X t_0.
\mathbf{lemma}\ sigma-finite-filtered-measure-natural-filtration:
 assumes exhausting-set: countable A (\bigcup A) = space M \land a. \ a \in A \Longrightarrow emeasure
M \ a \neq \infty \ \land a. \ a \in A \Longrightarrow \exists \ b \in borel. \ a = X \ t_0 - `b \cap space M
   shows sigma-finite-filtered-measure M (natural-filtration M t_0 X) t_0
proof (unfold-locales)
 have A \subseteq sets (restr-to-subalg M (natural-filtration M t_0 X t_0)) using exhaust-
ing\text{-}set by (simp\ add:\ sets\text{-}restr\text{-}to\text{-}subalg[OF\ subalgebra\text{-}natural\text{-}filtration}]\ sets\text{-}natural\text{-}filtration})
fast
 moreover have \bigcup A = space (restr-to-subalg M (natural-filtration M <math>t_0 X t_0))
unfolding space-restr-to-subalg using exhausting-set by simp
```

 $t_0)$ )  $a \neq \infty$  using calculation(1) exhausting-set(3)

**moreover have**  $\forall a \in A$ . emeasure (restr-to-subalg M (natural-filtration M  $t_0$  X

```
by (auto simp add: sets-restr-to-subalg[OF subalgebra-natural-filtration] emea-
sure-restr-to-subalg[OF\ subalgebra-natural-filtration])
 ultimately show \exists A. countable A \land A \subseteq sets (restr-to-subalg M (natural-filtration
M \ t_0 \ X \ t_0) \land \bigcup A = space \ (restr-to-subalq \ M \ (natural-filtration \ M \ t_0 \ X \ t_0)) \land
(\forall a \in A. \ emeasure \ (restr-to-subalg \ M \ (natural-filtration \ M \ t_0 \ X \ t_0)) \ a \neq \infty) using
exhausting-set by blast
 show \bigwedge i j. \llbracket t_0 \leq i; i \leq j \rrbracket \Longrightarrow sets (natural-filtration <math>M \ t_0 \ X \ i) \subseteq sets (natural-filtration M \ t_0 \ X \ i)
M t_0 X j) using filtered-measure.subalgebra-F[OF filtered-measure-natural-filtration]
by (simp add: subalgebra-def)
qed (auto intro: subalgebra-natural-filtration)
lemma finite-filtered-measure-natural-filtration:
 assumes finite-measure M
 shows finite-filtered-measure M (natural-filtration M t_0 X) t_0
 using finite-measure.axioms[OF assms] filtered-measure-natural-filtration by in-
tro-locales
end
Filtration generated by independent variables.
lemma (in prob-space) indep-set-natural-filtration:
 assumes t_0 \leq s \; s < t \; indep-vars \; (\lambda -. \; borel) \; X \; \{t_0..\}
  shows indep-set (natural-filtration M t_0 X s) (vimage-algebra (space M) (X t)
borel)
proof
 have indep-sets (\lambda i. \{Xi - `A \cap space\ M \mid A.\ A \in sets\ borel\}) (U) (range (case-bool
\{t_0..s\}\ \{t\})))
   using assms
  by (intro assms(3)[unfolded indep-vars-def, THEN conjunct2, THEN indep-sets-mono])
(auto simp add: case-bool-if)
  thus ?thesis unfolding indep-set-def using assms
   by (intro indep-sets-cong THEN iffD1, OF reft - indep-sets-collect-sigma of \lambda i.
\{X \ i - A \cap space \ M \mid A. \ A \in borel\} \ case-bool \ \{t_0...s\} \ \{t\}\}\}
       (simp add: sets-natural-filtration sets-vimage-algebra split: bool.split, simp,
intro Int-stableI, clarsimp, metis sets.Int vimage-Int Int-commute Int-left-absorb
Int-left-commute, force simp add: disjoint-family-on-def split: bool.split)
qed
```

# 6.2 Adapted Process

We call a collection a stochastic process X adapted if X i is F i-borel-measurable for all indices i.

```
locale adapted-process = filtered-measure M F t_0 for M F t_0 and X :: - \Rightarrow - \Rightarrow -:: \{second\text{-}countable\text{-}topology, banach}\} + assumes adapted[measurable]: \bigwedge i. \ t_0 \leq i \Longrightarrow X \ i \in borel\text{-}measurable \ (F i) begin
```

**lemma** adaptedE[elim]:

```
assumes \llbracket \bigwedge j \ i. \ t_0 \leq j \Longrightarrow j \leq i \Longrightarrow X \ j \in borel\text{-}measurable (F i) \rrbracket \Longrightarrow P
 shows P
 using assms using adapted by (metis dual-order.trans borel-measurable-subalgebra
sets-F-mono space-F)
lemma adaptedD:
 assumes t_0 \leq j j \leq i
 shows X j \in borel-measurable (F i) using assms adapted by meson
end
locale nat-adapted-process = adapted-process M F \theta :: nat X  for M F X
locale \ real-adapted-process = adapted-process M F 0 :: real \ X \ for \ M \ F \ X
\mathbf{sublocale}\ \mathit{nat-adapted-process} \subseteq \mathit{nat-filtered-measure}\ ..
sublocale real-adapted-process \subseteq real-filtered-measure ..
lemma (in filtered-measure) adapted-process-const-fun:
 assumes f \in borel-measurable (F t_0)
 shows adapted-process M F t_0 (\lambda - f)
 using measurable-from-subalg subalgebra-F assms by (unfold-locales) blast
lemma (in filtered-measure) adapted-process-const:
  shows adapted-process M F t_0 (\lambda i - c i) by (unfold-locales) simp
Again, we cover basic operations.
context adapted-process
begin
lemma compose-adapted:
 assumes \bigwedge i. t_0 \leq i \Longrightarrow f i \in borel-measurable borel
 shows adapted-process M F t_0 (\lambda i \xi. (f i) (X i \xi))
 by (unfold-locales) (intro measurable-compose[OF adapted assms])
lemma norm-adapted: adapted-process M F t_0 (\lambda i \xi. norm (X i \xi)) by (fastforce
intro: compose-adapted)
lemma scaleR-right-adapted:
 assumes adapted-process M F t_0 R
 shows adapted-process M F t_0 (\lambda i \xi. (R i \xi) *_R (X i \xi))
 using adapted-process.adapted[OF assms] adapted by (unfold-locales) simp
{\bf lemma}\ scale R\hbox{-}right\hbox{-}const\hbox{-}fun\hbox{-}adapted\colon
 assumes f \in borel-measurable (F t_0)
 shows adapted-process M F t_0 (\lambda i \xi. f \xi *_R (X i \xi))
 using assms by (fast intro: scaleR-right-adapted adapted-process-const-fun)
lemma scaleR-right-const-adapted: adapted-process M F t_0 (\lambda i \xi. c i *_R (X i \xi))
by (unfold-locales) simp
```

```
lemma add-adapted:
  assumes adapted-process M F t_0 Y
  shows adapted-process M F t_0 (\lambda i \xi. X i \xi + Y i \xi)
  using adapted-process.adapted[OF assms] adapted by (unfold-locales) simp
lemma diff-adapted:
  assumes adapted-process M F t_0 Y
  shows adapted-process M F t_0 (\lambda i \xi. X i \xi - Y i \xi)
  using adapted-process.adapted[OF assms] adapted by (unfold-locales) simp
lemma uminus-adapted: adapted-process MFt_0(-X) using scaleR-right-const-adapted of
\lambda-. -1] by (simp add: fun-Compl-def)
lemma partial-sum-adapted: adapted-process M F t_0 (\lambda n \xi. \sum i \in \{t_0..n\}. X i \xi)
proof (unfold-locales)
 fix i :: 'b
 have X j \in borel-measurable (F i) if t_0 \leq j j \leq i for j using that adapted by
  thus (\lambda \xi. \sum i \in \{t_0..i\}. X i \xi) \in borel-measurable (F i) by simp
qed
lemma partial-sum'-adapted: adapted-process M F t_0 (\lambda n \xi. \sum i \in \{t_0... < n\}. X i \xi)
proof (unfold-locales)
  \mathbf{fix}\ i::\ 'b
 have X j \in borel-measurable (F i) if t_0 \leq j j < i for j using that adapted by
 thus (\lambda \xi. \sum i \in \{t_0... < i\}. X i \xi) \in borel-measurable (F i) by simp
qed
end
In the dicrete time case, we have the following lemma which will be useful
later on.
\mathbf{lemma} \ (\mathbf{in} \ nat\text{-}adapted\text{-}process) \ partial\text{-}sum\text{-}Suc\text{-}adapted\text{:} \ nat\text{-}adapted\text{-}process \ M
F(\lambda n \xi. \sum i < n. X(Suc i) \xi)
\mathbf{proof} \ (\mathit{unfold-locales})
  \mathbf{fix} i
 have X j \in borel-measurable (F i) if j \leq i for j using that adapted D by blast
 thus (\lambda \xi. \sum i < i. X (Suc i) \xi) \in borel-measurable (F i) by auto
qed
lemma (in filtered-measure) adapted-process-sum:
  assumes \bigwedge i. i \in I \Longrightarrow adapted-process M F t_0 (X i)
  shows adapted-process M F t_0 (\lambda k \xi. \sum i \in I. X i k \xi)
proof -
  {
   fix i k assume i \in I and asm: t_0 \le k
```

```
then interpret adapted-process M F t_0 X i using assms by simp
    have X \ i \ k \in borel-measurable M \ X \ i \ k \in borel-measurable (F \ k) using mea-
surable-from-subalg subalgebras adapted asm by (blast, simp)
 thus ?thesis by (unfold-locales) simp
qed
An adapted process is necessarily a stochastic process.
\mathbf{sublocale}\ adapted-process \subseteq stochastic-process \mathbf{using}\ measurable-from-subalg sub-
algebras adapted by (unfold-locales) blast
\mathbf{sublocale}\ \mathit{nat-adapted-process} \subseteq \mathit{nat-stochastic-process}\ ..
\mathbf{sublocale}\ \mathit{real-adapted-process} \subseteq \mathit{real-stochastic-process}\ ..
A stochastic process is always adapted to the natural filtration it generates.
{f lemma} (in stochastic-process) adapted-process-natural-filtration: <math>adapted-process
M (natural-filtration M t_0 X) t_0 X
 using filtered-measure-natural-filtration
 by (intro-locales) (auto simp add: natural-filtration-def intro!: adapted-process-axioms.intro
measurable-family-vimage-algebra)
6.3
       Progressively Measurable Process
locale progressive-process = filtered-measure M F t_0 for M F t_0 and X :: - \Rightarrow -
\Rightarrow - :: { second-countable-topology, banach} +
assumes progressive[measurable]: \bigwedge t. t_0 \le t \Longrightarrow (\lambda(i, x). X i x) \in borel-measurable
(restrict-space borel \{t_0..t\} \bigotimes_M F t)
begin
lemma progressiveD:
 assumes S \in borel
 shows (\lambda(j, \xi). X j \xi) - 'S \cap (\{t_0..i\} \times space M) \in (restrict\text{-}space borel \{t_0..i\}
\bigotimes_M F(i)
 using measurable-sets[OF progressive, OF - assms, of i]
 by (cases t_0 \le i) (auto simp add: space-restrict-space sets-pair-measure space-pair-measure)
end
locale nat-progressive-process = progressive-process M F 0 :: nat X  for M F X
locale real-progressive-process = progressive-process M F 0 :: real X \text{ for } M F X
lemma (in filtered-measure) progressive-process-const-fun:
 assumes f \in borel-measurable (F t_0)
 shows progressive-process M F t_0 (\lambda - f)
proof (unfold-locales)
  fix i assume asm: t_0 \leq i
 have f \in borel-measurable (F i) using borel-measurable-mono[OF order.refl asm]
assms by blast
```

```
thus case-prod (\lambda-. f) \in borel-measurable (restrict-space borel \{t_0...i\} \bigotimes_M Fi)
using measurable-compose[OF measurable-snd] by simp
qed
lemma (in filtered-measure) progressive-process-const:
 assumes c \in borel-measurable borel
 shows progressive-process M F t_0 (\lambda i -. c i)
  using assms by (unfold-locales) (auto simp add: measurable-split-conv intro!:
measurable-compose[OF measurable-fst] measurable-restrict-space1)
{\bf context}\ progressive\text{-}process
begin
lemma compose-progressive:
 assumes case-prod f \in borel-measurable borel
 shows progressive-process M F t_0 (\lambda i \xi. (f i) (X i \xi))
proof
 fix i assume asm: t_0 \leq i
 have (\lambda(j, \xi). (j, X j \xi)) \in (\textit{restrict-space borel } \{t_0..i\} \bigotimes_M F i) \rightarrow_M \textit{borel } \bigotimes_M F i)
    using progressive [OF asm] measurable-fst" [OF measurable-restrict-space1, OF
measurable-id
   by (auto simp add: measurable-pair-iff measurable-split-conv)
 moreover have (\lambda(j, \xi). f j (X j \xi)) = case-prod f o ((\lambda(j, y). (j, y)) o (\lambda(j, \xi).
(j, X j \xi)) by fastforce
  ultimately show (\lambda(j, \xi), (fj), (Xj\xi)) \in borel-measurable (restrict-space borel)
\{t_0...i\} \bigotimes_M F(i) using assms by (simp add: borel-prod)
lemma norm-progressive: progressive-process M F t_0 (\lambda i \xi. norm (X i \xi)) us-
ing measurable-compose OF progressive borel-measurable-norm by (unfold-locales)
simp
{f lemma}\ scale R	ext{-}right	ext{-}progressive:
 assumes progressive-process M F t_0 R
 shows progressive-process M F t_0 (\lambda i \xi. (R i \xi) *_R (X i \xi))
  using progressive-process.progressive[OF assms] by (unfold-locales) (simp add:
progressive assms)
lemma scaleR-right-const-fun-progressive:
 assumes f \in borel-measurable (F t_0)
 shows progressive-process M F t_0 (\lambda i \xi. f \xi *_R (X i \xi))
 using assms by (fast intro: scaleR-right-progressive progressive-process-const-fun)
{\bf lemma}\ scale R\hbox{-}right\hbox{-}const\hbox{-}progressive\hbox{:}
 assumes c \in borel-measurable borel
 shows progressive-process M F t_0 (\lambda i \xi. c i *_R (X i \xi))
 using assms by (fastforce intro: scaleR-right-progressive progressive-process-const)
```

```
lemma add-progressive:
 assumes progressive-process M F t_0 Y
 shows progressive-process M F t_0 (\lambda i \xi. X i \xi + Y i \xi)
  using progressive-process.progressive[OF assms] by (unfold-locales) (simp add:
progressive assms)
lemma diff-progressive:
 assumes progressive-process M F t_0 Y
 shows progressive-process M F t_0 (\lambda i \xi. X i \xi - Y i \xi)
  using progressive-process.progressive[OF assms] by (unfold-locales) (simp add:
progressive assms)
lemma uminus-progressive: progressive-process M F t_0 (-X) using scaleR-right-const-progressive[of
\lambda-. -1] by (simp add: fun-Compl-def)
end
A progressively measurable process is also adapted.
sublocale progressive-process \subseteq adapted-process using measurable-compose-rev[OF]
progressive measurable-Pair1 |
 unfolding prod.case space-restrict-space
 by unfold-locales simp
sublocale nat-progressive-process \subseteq nat-adapted-process ...
\mathbf{sublocale}\ real	ext{-}progressive	ext{-}process \subseteq real	ext{-}adapted	ext{-}process ..
In the discrete setting, adaptedness is equivalent to progressive measurabil-
ity.
theorem nat-progressive-iff-adapted: nat-progressive-process MFX \longleftrightarrow nat-adapted-process
MFX
proof (intro iffI)
 assume asm: nat\text{-}progressive\text{-}process \ M \ F \ X
 interpret nat-progressive-process M F X by (rule asm)
 show nat-adapted-process M F X ..
next
 assume asm: nat-adapted-process MFX
 interpret nat-adapted-process M F X by (rule asm)
 show nat-progressive-process M F X
 proof (unfold-locales, intro borel-measurableI)
   fix S :: 'b \ set \ and \ i :: nat \ assume \ open-S: open \ S
     fix j assume asm: j < i
    hence X j - S \cap Space M \in F i using adaptedD[of j, THEN measurable-sets]
space-F open-S by fastforce
     moreover have case-prod X - S \cap \{j\} \times space M = \{j\} \times (X j - S \cap S)
space M) for j by fast
     moreover have \{j :: nat\} \in restrict\text{-space borel } \{0..i\} \text{ using } asm \text{ by } (simp)
add: sets-restrict-space-iff)
```

```
ultimately have case-prod X – 'S \cap \{j\} × space M \in restrict-space borel \{0..i\} \bigotimes_M F i by simp
}
hence (\lambda j. (\lambda(x, y). \ X \ x \ y) – 'S \cap \{j\} × space M) '\{..i\} \subseteq restrict-space borel \{0..i\} \bigotimes_M F i by blast
moreover have case-prod X – 'S \cap space (restrict-space borel \{0..i\} \bigotimes_M F i) = (\bigcup_j \leq i. case-prod X – 'S \cap \{j\} × space M) unfolding space-pair-measure space-restrict-space space-F by force
ultimately show case-prod X – 'S \cap space (restrict-space borel \{0..i\} \bigotimes_M F i) \in restrict-space borel \{0..i\} \bigotimes_M F i by (metis sets.countable-UN) qed qed
```

### 6.4 Predictable Process

We introduce the constant  $\Sigma_P$  to denote the predictable  $\sigma$ -algebra.

**context** linearly-filtered-measure **begin** 

 ${\bf lemma}\ measurable ext{-}predictable ext{-}sigma ext{-}snd:$ 

```
definition \Sigma_P :: ('b \times 'a) measure where predictable-sigma: \Sigma_P \equiv sigma \ (\{t_0..\} \times space \ M) \ (\{\{s<..t\} \times A \mid A \ s \ t. \ A \in F \ s \land t_0 \leq s \land s < t\} \cup \{\{t_0\} \times A \mid A. \ A \in F \ t_0\})
```

lemma space-predictable-sigma[simp]: space  $\Sigma_P = (\{t_0..\} \times space\ M)$  unfolding predictable-sigma space-measure-of-conv by blast

```
lemma sets-predictable-sigma: sets \Sigma_P = sigma\text{-sets} \ (\{t_0..\} \times space \ M) \ (\{\{s<..t\} \times A \mid A \ s \ t. \ A \in F \ s \land t_0 \leq s \land s < t\} \cup \{\{t_0\} \times A \mid A. \ A \in F \ t_0\}) unfolding predictable-sigma using space-F sets.sets-into-space by (subst sets-measure-of) fastforce+
```

```
assumes countable \mathcal{I} \mathcal{I} \subseteq \{\{s<..t\} \mid s \ t. \ t_0 \leq s \land s < t\} \ \{t_0<..\} \subseteq (\bigcup \mathcal{I})
  shows snd \in \Sigma_P \to_M F t_0
proof (intro measurableI)
  fix S :: 'a \text{ set assume } asm: S \in F t_0
  have countable: countable ((\lambda I.\ I \times S)\ '\mathcal{I}) using assms(1) by blast
  have (\lambda I. \ I \times S) '\mathcal{I} \subseteq \{\{s < ...t\} \times A \mid A \ s \ t. \ A \in F \ s \land t_0 \leq s \land s < t\} using
sets-F-mono[OF order-reft, THEN subsetD, OF - asm] assms(2) by blast
 hence (\bigcup I \in \mathcal{I}. \ I \times S) \cup \{t_0\} \times S \in \Sigma_P  unfolding sets-predictable-sigma using
asm\ \mathbf{by}\ (intro\ sigma-sets-UniON[OF\ countable]\ sigma-sets.Basic]
sigma-sets.Basic) blast+
 moreover have snd - S \cap space \Sigma_P = \{t_0..\} \times S \text{ using } sets.sets-into-space [OF] \}
asm] by fastforce
  moreover have \{t_0\} \cup \{t_0 < ..\} = \{t_0..\} by auto
  moreover have (\bigcup I \in \mathcal{I}. \ I \times S) \cup \{t_0\} \times S = \{t_0..\} \times S \text{ using } assms(2,3)
calculation(3) by fastforce
  ultimately show snd - Snace \Sigma_P \in \Sigma_P by argo
qed (auto)
```

```
{\bf lemma}\ measurable\text{-}predictable\text{-}sigma\text{-}fst\text{:}
  assumes countable \mathcal{I} \mathcal{I} \subseteq \{\{s<..t\} \mid s \ t. \ t_0 \leq s \land s < t\} \ \{t_0<..\} \subseteq (\bigcup \mathcal{I})
  shows fst \in \Sigma_P \to_M borel
proof -
  have A \times space M \in sets \Sigma_P if A \in sigma-sets \{t_0..\} \{\{s<..t\} \mid s \ t. \ t_0 \leq s \land s\}
\langle t \rangle for A unfolding sets-predictable-sigma using that
  proof (induction rule: sigma-sets.induct)
   case (Basic\ a)
   thus ?case using space-F sets.top by blast
  next
   have (\{t_0..\} - a) \times space M = \{t_0..\} \times space M - a \times space M by blast
   then show ?case using Compl(2)[THEN sigma-sets.Compl] by presburger
  next
    case (Union a)
   have \bigcup (range a) \times space M = \bigcup (range (\lambda i.\ a\ i \times space\ M)) by blast
   then show ?case using Union(2)[THEN sigma-sets.Union] by presburger
  moreover have restrict-space borel \{t_0..\} = sigma\ \{t_0..\}\ \{\{s<..t\} \mid s\ t.\ t_0 \le s
\land s < t
  proof -
    have sigma-sets \{t_0..\} ((\cap)\ \{t_0..\}\ 'sigma-sets\ UNIV\ (range\ greaterThan)) =
sigma-sets \{t_0..\} \{\{s<..t\} | s \ t. \ t_0 \le s \land s < t\}
   proof (intro sigma-sets-eqI ; clarify)
      fix A :: 'b \text{ set assume } asm: A \in sigma-sets UNIV (range greaterThan)
      thus \{t_0..\} \cap A \in sigma\text{-sets } \{t_0..\} \{\{s<..t\} \mid s \ t. \ t_0 \leq s \land s < t\}
      proof (induction rule: sigma-sets.induct)
       case (Basic\ a)
       then obtain s where s: a = \{s < ...\} by blast
       show ?case
       proof (cases t_0 \leq s)
         case True
         hence *: \{t_0..\} \cap a = (\bigcup i \in \mathcal{I}. \{s<..\} \cap i) using s \ assms(3) by force
         have ((\cap) \{s<...\} '\mathcal{I}) \subseteq sigma-sets \{t_0...\} \{\{s<..t\} \mid s \ t. \ t_0 \leq s \land s < t\}
         proof (clarify)
           fix A assume A \in \mathcal{I}
          then obtain s' t' where A: A = \{s' < ...t'\} t_0 \le s' s' < t' using assms(2)
by blast
           hence \{s<...\} \cap A = \{max \ s \ s'<...t'\} by fastforce
            moreover have t_0 \leq max \ s \ s' using A True by linarith
            moreover have \max s s' < t' if s < t' using A that by linarith
           moreover have \{s<...\} \cap A = \{\} if \neg s < t' using A that by force
            ultimately show \{s<...\} \cap A \in sigma-sets \{t_0...\} \{\{s<...t\} \mid s \ t. \ t_0 \leq s
\land s < t} by (cases s < t') (blast, simp add: sigma-sets.Empty)
          thus ?thesis unfolding * using assms(1) by (intro sigma-sets-UNION)
auto
       next
```

```
case False
         hence \{t_0..\} \cap a = \{t_0..\} using s by force
         thus ?thesis using sigma-sets-top by auto
       qed
     next
       case (Compl\ a)
       have \{t_0..\} \cap (UNIV - a) = \{t_0..\} - (\{t_0..\} \cap a) by blast
       then show ?case using Compl(2)[THEN\ sigma-sets.Compl] by presburger
     next
       case (Union \ a)
       have \{t_0..\} \cap \bigcup (range \ a) = \bigcup (range \ (\lambda i. \ \{t_0..\} \cap a \ i)) by blast
       then show ?case using Union(2)[THEN \ sigma-sets.Union] by presburger
     qed (simp add: sigma-sets.Empty)
   next
     fix s t assume asm: t_0 \le s \ s < t
     hence *: \{s<...t\} = \{s<...\} \cap (\{t_0..\} - \{t<...\}) by force
    have \{s<...\} \in sigma-sets \{t_0...\} ((\cap) \{t_0...\} 'sigma-sets UNIV (range greaterThan))
using asm by (intro sigma-sets.Basic) auto
      moreover have \{t_0..\} - \{t<..\} \in sigma-sets \{t_0..\} ((\cap) \{t_0..\} 'sigma-sets \}
UNIV (range greaterThan)) using asm by (intro sigma-sets.Compl sigma-sets.Basic)
     ultimately show \{s<..t\} \in sigma-sets \{t_0..\} ((\cap) \{t_0..\} \ `sigma-sets \ UNIV
(range\ greaterThan))\ \mathbf{unfolding}*Int-range-binary[of\ \{s<..\}]\ \mathbf{by}\ (intro\ sigma-sets-Inter[OF\ sets])
- binary-in-sigma-sets]) auto
   qed
    thus ?thesis unfolding borel-Ioi restrict-space-def emeasure-sigma by (force
intro: sigma-eqI)
 ged
 ultimately have restrict-space borel \{t_0..\} \bigotimes_M sigma (space M) \{\} \subseteq sets \Sigma_P
   unfolding sets-pair-measure space-restrict-space space-measure-of-conv
   using space-predictable-sigma sets.sigma-algebra-axioms[of \Sigma_P]
   by (intro sigma-algebra.sigma-sets-subset) (auto simp add: sigma-sets-empty-eq
sets-measure-of-conv)
  moreover have space (restrict-space borel \{t_0..\} \bigotimes_M sigma (space M) \{\}) =
space \Sigma_P by (simp add: space-pair-measure)
  moreover have fst \in restrict\text{-}space \ borel\ \{t_0..\}\ \bigotimes_{M} \ sigma\ (space\ M)\ \{\} \rightarrow_{M}
borel by (fastforce intro: measurable-fst" [OF measurable-restrict-space1, of \lambda x. x])
  ultimately show ?thesis by (meson borel-measurable-subalgebra)
qed
end
locale predictable\text{-}process = linearly\text{-}filtered\text{-}measure } M \ F \ t_0 \ \text{for } M \ F \ t_0 \ \text{and } X ::
- \Rightarrow - \Rightarrow - :: \{second\text{-}countable\text{-}topology, banach} +
 assumes predictable: (\lambda(t, x), X t x) \in borel-measurable \Sigma_P
begin
```

```
\label{eq:continuous} \textbf{lemmas} \ predictable D = measurable-sets [OF \ predictable, \ unfolded \ space-predictable-sigma] \textbf{end}
```

```
locale \ nat-predictable-process = predictable-process \ M \ F \ 0 :: nat \ X \ for \ M \ F \ X
locale real-predictable-process = predictable-process M F 0 :: real X for M F X
\mathbf{lemma} \ (\mathbf{in} \ \mathit{nat-filtered-measure}) \ \mathit{measurable-predictable-sigma-snd'} :
 shows snd \in \Sigma_P \to_M F \theta
 by (intro measurable-predictable-sigma-snd[of range (\lambda x. {Suc x})]) (force | simp
add: qreaterThan-\theta)+
lemma (in nat-filtered-measure) measurable-predictable-sigma-fst':
 shows fst \in \Sigma_P \to_M borel
 by (intro measurable-predictable-sigma-fst[of range (\lambda x. \{Suc\ x\})]) (force | simp
add: qreaterThan-\theta)+
lemma (in real-filtered-measure) measurable-predictable-sigma-snd':
 shows snd \in \Sigma_P \to_M F \theta
 using real-arch-simple by (intro measurable-predictable-sigma-snd[of range (\lambda x::nat.
\{0 < ... real (Suc x)\}\} (fastforce intro: add-increasing)+
lemma (in real-filtered-measure) measurable-predictable-sigma-fst':
 shows fst \in \Sigma_P \to_M borel
 using real-arch-simple by (intro measurable-predictable-sigma-fst of range (\lambda x:: nat.
\{0 < ... real (Suc x)\}\} (fastforce intro: add-increasing)+
We show sufficient conditions for functions constant in one argument to
constitute a predictable process. In contrast to the cases before, this is not
a triviality.
lemma (in linearly-filtered-measure) predictable-process-const-fun:
 assumes snd \in \Sigma_P \to_M F t_0 f \in borel\text{-}measurable (F t_0)
   shows predictable-process M F t_0 (\lambda -... f)
  using measurable-compose-rev[OF\ assms(2)]\ assms(1) by (unfold-locales) (auto
simp add: measurable-split-conv)
\mathbf{lemma} \ (\mathbf{in} \ \mathit{nat-filtered-measure}) \ \mathit{predictable-process-const-fun'[intro]} :
 assumes f \in borel-measurable (F \ \theta)
 shows nat-predictable-process M F (\lambda-. f)
 using assms by (intro predictable-process-const-fun OF measurable-predictable-sigma-snd',
THEN nat-predictable-process.intro])
lemma (in real-filtered-measure) predictable-process-const-fun'[intro]:
  assumes f \in borel-measurable (F \ \theta)
 shows real-predictable-process M F (\lambda - f)
 using assms by (intro predictable-process-const-fun OF measurable-predictable-sigma-snd',
THEN real-predictable-process.intro])
```

```
lemma (in linearly-filtered-measure) predictable-process-const:
 assumes fst \in borel-measurable \Sigma_P c \in borel-measurable borel
 shows predictable-process M F t_0 (\lambda i -. c i)
 using assms by (unfold-locales) (simp add: measurable-split-conv)
\mathbf{lemma} \ (\mathbf{in} \ \mathit{linearly-filtered-measure}) \ \mathit{predictable-process-const-const}[\mathit{intro}] :
  shows predictable-process M F t_0 (\lambda - c)
 by (unfold-locales) simp
lemma (in nat-filtered-measure) predictable-process-const'[intro]:
 assumes c \in borel-measurable borel
 shows nat-predictable-process M F (\lambda i - c i)
 using assms by (intro predictable-process-const[OF measurable-predictable-sigma-fst',
THEN nat-predictable-process.intro])
lemma (in real-filtered-measure) predictable-process-const'[intro]:
 assumes c \in borel-measurable borel
 shows real-predictable-process M F (\lambda i - c i)
 using assms by (intro predictable-process-const[OF measurable-predictable-siqma-fst',
THEN real-predictable-process.intro])
context predictable-process
begin
lemma compose-predictable:
  assumes fst \in borel-measurable \Sigma_P case-prod f \in borel-measurable borel
 shows predictable-process M F t_0 (\lambda i \xi. (f i) (X i \xi))
proof
 have (\lambda(i, \xi), (i, X i \xi)) \in \Sigma_P \to_M borel \bigotimes_M borel using predictable assms(1)
by (auto simp add: measurable-pair-iff measurable-split-conv)
  moreover have (\lambda(i, \xi), f(i, X(i, \xi))) = case-prod f(i, \xi), (i, X(i, \xi)) by
fastforce
 ultimately show (\lambda(i, \xi). fi(Xi\xi)) \in borel-measurable \Sigma_P unfolding borel-prod
using assms by simp
qed
lemma norm-predictable: predictable-process M F t_0 (\lambda i \ \xi. norm (X i \ \xi)) using
measurable-compose[OF\ predictable\ borel-measurable-norm]
 by (unfold-locales) (simp add: prod.case-distrib)
lemma scaleR-right-predictable:
 assumes predictable-process M F t_0 R
 shows predictable-process M F t_0 (\lambda i \xi. (R i \xi) *_R (X i \xi))
  using predictable predictable-process.predictable[OF assms] by (unfold-locales)
(auto simp add: measurable-split-conv)
\mathbf{lemma}\ scaleR-right-const-fun-predictable:
 assumes snd \in \Sigma_P \to_M F t_0 f \in borel\text{-}measurable (F t_0)
```

```
shows predictable-process M F t_0 (\lambda i \xi. f \xi *_R (X i \xi))
 using assms by (fast intro: scaleR-right-predictable predictable-process-const-fun)
\mathbf{lemma}\ scaleR-right-const-predictable:
  assumes fst \in borel-measurable \Sigma_P c \in borel-measurable borel
 shows predictable-process M F t_0 (\lambda i \xi. c i *_R (X i \xi))
 using assms by (fastforce intro: scaleR-right-predictable predictable-process-const)
lemma scale
R-right-const'-predictable: predictable-process M F t_0 (\lambda i \xi. c*_R (X i
  by (fastforce intro: scaleR-right-predictable)
lemma add-predictable:
  assumes predictable-process M F t_0 Y
  shows predictable-process M F t_0 (\lambda i \xi. X i \xi + Y i \xi)
  using predictable predictable-process.predictable[OF assms] by (unfold-locales)
(auto simp add: measurable-split-conv)
lemma diff-predictable:
  assumes predictable-process M F t_0 Y
 shows predictable-process M F t_0 (\lambda i \ \xi. X i \ \xi - Y \ i \ \xi)
  using predictable predictable-process.predictable[OF assms] by (unfold-locales)
(auto simp add: measurable-split-conv)
lemma uminus-predictable: predictable-process MFt_0(-X) using scaleR-right-const'-predictable[of
-1] by (simp add: fun-Compl-def)
end
Every predictable process is also progressively measurable.
sublocale predictable-process \subseteq progressive-process
proof (unfold-locales)
  fix i :: 'b assume asm: t_0 \leq i
   \mathbf{fix}\ S :: ('b \times 'a)\ set\ \mathbf{assume}\ S \in \{\{s{<}..t\} \times A \mid A\ s\ t.\ A \in F\ s\ \land\ t_0 \leq s\ \land\ s
< t \} \cup \{ \{t_0\} \times A \mid A. A \in F t_0 \}
   hence (\lambda x.\ x) - 'S \cap (\{t_0...i\} \times space\ M) \in restrict\text{-space borel } \{t_0...i\} \bigotimes_M F
   proof
     assume S \in \{\{s < ...t\} \times A \mid A \ s \ t. \ A \in F \ s \land t_0 \le s \land s < t\}
      then obtain s \ t \ A where S-is: S = \{s < ...t\} \times A \ t_0 \le s \ s < t \ A \in F \ s by
blast
       hence (\lambda x. \ x) - S \cap (\{t_0...i\} \times space \ M) = \{s < ... min \ i \ t\} \times A \ using
sets.sets-into-space[\mathit{OF}\ S-is(4)] by auto
     then show ?thesis using S-is sets-F-mono[of s i] by (cases s \leq i) (fastforce
simp add: sets-restrict-space-iff)+
     assume S \in \{\{t_0\} \times A \mid A. A \in F t_0\}
```

then obtain A where S-is:  $S = \{t_0\} \times A \ A \in F \ t_0 \ \text{by} \ blast$ 

```
hence (\lambda x. x) - S \cap (\{t_0...i\} \times space M) = \{t_0\} \times A using asm sets.sets-into-space OF
S-is(2)] by auto
     thus ?thesis using S-is(2) sets-F-mono[OF order-refl asm] asm by (fastforce
simp add: sets-restrict-space-iff)
   ged
   hence (\lambda x. x) - 'S \cap space (restrict-space borel \{t_0..i\} \bigotimes_M Fi) \in restrict-space
borel \{t_0..i\} \bigotimes_M F i by (simp \ add: space-pair-measure \ space-F[OF \ asm])
 moreover have \{\{s<..t\} \times A \mid A \text{ s t. } A \in sets (F \text{ s}) \land t_0 \leq s \land s < t\} \cup \{\{t_0\}\}\}
\times A |A. A \in sets (F t<sub>0</sub>)} \subseteq Pow ({t<sub>0</sub>..} \times space M) using sets.sets-into-space by
force
  ultimately have (\lambda x. \ x) \in restrict\text{-space borel} \ \{t_0..i\} \bigotimes_M F \ i \to_M \Sigma_P \text{ us-}
ing space-F[OF asm] by (intro measurable-sigma-sets[OF sets-predictable-sigma])
(fast, force simp add: space-pair-measure)
 thus case-prod X \in borel-measurable (restrict-space borel \{t_0...i\} \bigotimes_M F(i)) using
predictable by simp
qed
sublocale nat-predictable-process \subseteq nat-progressive-process ...
sublocale real-predictable-process \subseteq real-progressive-process ...
The following lemma characterizes predictability in a discrete-time setting.
lemma (in nat-filtered-measure) sets-in-filtration:
  assumes (\bigcup i. \{i\} \times A \ i) \in \Sigma_P
  shows A (Suc i) \in F i A 0 \in F 0
  using assms unfolding sets-predictable-sigma
proof (induction (\bigcup i. \{i\} \times A \ i) arbitrary: A)
  {f case}\ Basic
  {
   assume \exists S. (\bigcup i. \{i\} \times A \ i) = \{0\} \times S
   then obtain S where S: (\bigcup i. \{i\} \times A \ i) = \{bot\} \times S unfolding bot-nat-def
   hence S \in F bot using Basic by (fastforce simp add: times-eq-iff bot-nat-def)
   moreover have A \ i = \{\} if i \neq bot for i using that S by blast
   moreover have A bot = S using S by blast
   ultimately have A (Suc i) \in F i A \theta \in F \theta for i unfolding bot-nat-def by
(auto simp add: bot-nat-def)
  }
 note * = this
   assume \nexists S. (\bigcup i. \{i\} \times A \ i) = \{0\} \times S
    then obtain s \ t \ B where B: (\bigcup i. \{i\} \times A \ i) = \{s < ... t\} \times B \ B \in sets \ (F \ s)
s < t using Basic by auto
   hence A \ i = B \ \text{if} \ i \in \{s < ... t\} \ \text{for} \ i \ \text{using} \ that \ \text{by} \ fast
   moreover have A i = \{\} if i \notin \{s < ... t\} for i using B that by fastforce
   ultimately have A(Suc\ i) \in F\ i\ A\ \theta \in F\ \theta for i unfolding bot-nat-def using
B sets-F-mono by (auto simp add: bot-nat-def) (metis less-Suc-eq-le sets.empty-sets
subset-eq)
  }
```

```
note ** = this
 show A (Suc i) \in sets (F i) A 0 \in sets (F 0) using *(1)[of i] *(2) **(1)[of i]
**(2) by blast+
\mathbf{next}
 case Empty
   case 1
   then show ?case using Empty by simp
  next
   then show ?case using Empty by simp
 }
next
 case (Compl a)
 have a-in: a \subseteq \{0..\} \times space \ M \ using \ Compl(1) \ sets.sets-into-space \ sets-predictable-sigma
space-predictable-sigma by metis
 hence A-in: A i \subseteq space \ M for i \ using \ Compl(4) by blast
 have a: a = \{0..\} \times space \ M - (\bigcup i. \{i\} \times A \ i) \ using \ a-in \ Compl(4) \ by \ blast
 also have ... = -(\bigcap j. - (\{j\} \times (space M - A j))) by blast
 also have ... = (\bigcup j. \{j\} \times (space M - A j)) by blast
  finally have *: (space\ M-A\ (Suc\ i))\in F\ i\ (space\ M-A\ 0)\in F\ 0 using
Compl(2,3) by auto
 {
   case 1
     then show ?case using * A-in by (metis bot-nat-0.extremum double-diff
sets.Diff sets.top sets-F-mono sets-le-imp-space-le space-F)
next
   case 2
     then show ?case using * A-in by (metis bot-nat-0.extremum double-diff
sets.Diff sets.top sets-F-mono sets-le-imp-space-le space-F)
 }
next
 case (Union a)
  have a-in: a \ i \subseteq \{0..\} \times space \ M \ for \ i \ using \ Union(1) \ sets.sets-into-space
sets-predictable-sigma space-predictable-sigma by metis
 hence A-in: A i \subseteq space \ M for i \ using \ Union(4) by blast
  have snd \ x \in snd \ (a \ i \cap (\{fst \ x\} \times space \ M)) \ \text{if} \ x \in a \ i \ \text{for} \ i \ x \ \text{using} \ that
a-in by fastforce
 hence a-i: a i = (\bigcup j. \{j\} \times (snd \ (a \ i \cap (\{j\} \times space \ M)))) for i by force
  have A-i: A i = snd ' (\bigcup (range \ a) \cap (\{i\} \times space \ M)) for i unfolding
Union(4) using A-in by force
 have *: snd '(a j \cap (\{Suc\ i\} \times space\ M)) \in F \ i \ snd '(a j \cap (\{0\} \times space\ M))
\in F \ 0 \ \mathbf{for} \ j \ \mathbf{using} \ Union(2,3)[OF \ a-i] \ \mathbf{by} \ auto
 {
   case 1
   have (\bigcup j. \ snd \ (a \ j \cap (\{Suc \ i\} \times space \ M))) \in F \ i \ using * by \ fast
    moreover have (\bigcup j. \ snd \ (a \ j \cap (\{Suc \ i\} \times space \ M))) = snd \ (\bigcup \ (range)
a) \cap (\{Suc\ i\} \times space\ M)) by fast
   ultimately show ?case using A-i by metis
```

```
next
            case 2
            have (\bigcup j. \ snd \ `(a \ j \cap (\{0\} \times space \ M))) \in F \ 0 \ using * by \ fast
            moreover have (\bigcup j. \ snd \ (a \ j \cap (\{0\} \times space \ M))) = snd \ (\bigcup \ (range \ a) \cap \{0\} \cap \{0
(\{\theta\} \times space\ M)) by fast
            ultimately show ?case using A-i by metis
qed
This leads to the following useful fact.
lemma (in nat-predictable-process) adapted-Suc: nat-adapted-process M F (\lambda i. X)
(Suc\ i)
proof (unfold-locales, intro borel-measurableI)
      fix S :: 'b \ set \ and \ i \ assume \ open-S: \ open \ S
      have \{Suc\ i\} = \{i < ... Suc\ i\} by fastforce
        hence \{Suc\ i\} \times space\ M \in \Sigma_P \text{ using } space\text{-}F[symmetric,\ of\ i] \text{ unfolding}
sets-predictable-sigma by (intro sigma-sets.Basic) blast
        moreover have case-prod X -' S \cap (UNIV \times space M) \in \Sigma_P unfolding
atLeast-0[symmetric] using open-S by (intro predictableD, simp add: borel-open)
        ultimately have case-prod X - 'S \cap (\{Suc\ i\} \times space\ M) \in \Sigma_P unfolding
sets-predictable-sigma using space-F sets.sets-into-space
                  by (subst Times-Int-distrib1 of {Suc i} UNIV space M, simplified], subst
inf.commute, subst Int-assoc[symmetric], subst Int-range-binary)
                      (intro\ sigma-sets-Inter\ binary-in-sigma-sets,\ fast) +
     moreover have case-prod X - S \cap (\{Suc\ i\} \times space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (X\ (Suc\ i\} \times Space\ M) = \{Suc\ i\} \times (
i) - 'S \cap space M) by (auto simp add: le-Suc-eq)
     moreover have ... = (\bigcup j. \{j\} \times (if j = Suc \ i \ then \ (X \ (Suc \ i) - `S \cap space \ M)
else {})) by (force split: if-splits)
     ultimately have (\bigcup j. \{j\} \times (if j = Suc \ i \ then \ (X \ (Suc \ i) - `S \cap space \ M) \ else
\{\})) \in \Sigma_P by argo
       thus X (Suc i) – 'S \cap space (F i) \in sets (F i) using sets-in-filtration[of \lambda j.
if j = Suc \ i \ then \ (X \ (Suc \ i) - `S \cap space \ M) \ else \ \{\}] \ space-F[OF \ zero-le] \ by
presburger
qed
The following lemma characterizes predictability in the discrete setting.
theorem nat-predictable-process-iff: nat-predictable-process MFX \longleftrightarrow nat-adapted-process
M F (\lambda i. X (Suc i)) \wedge X \theta \in borel\text{-}measurable (F \theta)
proof (intro iffI)
     assume asm: nat-adapted-process M F (\lambda i. X (Suc i)) \wedge X \theta \in borel-measurable
(F \theta)
      interpret nat-adapted-process M F \lambda i. X (Suc i) using asm by blast
      have (\lambda(x, y), X x y) \in borel-measurable \Sigma_P
       proof (intro borel-measurableI)
            fix S :: 'b \ set \ assume \ open-S: \ open \ S
            have \{i\} \times (X \ i - S \cap space \ M) \in sets \ \Sigma_P \ for \ i
            proof (cases i)
                   case \theta
                   then show ?thesis unfolding sets-predictable-sigma
```

```
using measurable-sets[OF - borel-open[OF open-S], of X 0 F 0] asm by auto
   next
     case (Suc\ i)
     have \{Suc\ i\} = \{i < ... Suc\ i\} by fastforce
     then show ?thesis unfolding sets-predictable-sigma
       using measurable-sets[OF adapted borel-open[OF open-S], of i]
       by (intro sigma-sets.Basic, auto simp add: Suc)
   moreover have (\lambda(x, y). X x y) - S \cap Space \Sigma_P = (\bigcup i. \{i\} \times (X i - S \cap S))
space M)) by fastforce
   ultimately show (\lambda(x, y). X x y) - S \cap space \Sigma_P \in sets \Sigma_P by simp
 thus nat-predictable-process M F X by (unfold-locales)
next
  assume asm: nat\text{-}predictable\text{-}process M F X
 interpret nat-predictable-process M F X by (rule asm)
 show nat-adapted-process M F (\lambda i. X (Suc i)) \wedge X \theta \in borel-measurable (F \theta)
using adapted-Suc by simp
qed
end
theory Martingale
 imports Stochastic-Process Conditional-Expectation-Banach
begin
```

# 7 Martingales

The following locales are necessary for defining martingales.

#### 7.1 Additional Locale Definitions

```
 \begin{array}{l} \textbf{locale} \ sigma\text{-}finite\text{-}adapted\text{-}process = sigma\text{-}finite\text{-}filtered\text{-}measure} \ M\ F\ t_0\ +\ adapted\text{-}process \\ M\ F\ t_0\ X\ \textbf{for}\ M\ F\ t_0\ X \end{array}
```

```
locale nat-sigma-finite-adapted-process = sigma-finite-adapted-process M F 0 :: nat X for M F X locale real-sigma-finite-adapted-process = sigma-finite-adapted-process M F 0 :: real X for M F X sublocale nat-sigma-finite-adapted-process \subseteq nat-sigma-finite-filtered-measure \dots sublocale real-sigma-finite-adapted-process \subseteq real-sigma-finite-filtered-measure \dots locale finite-adapted-process = finite-filtered-measure M F t_0 + adapted-process M F t_0 X for M F t_0 X sublocale finite-adapted-process \subseteq sigma-finite-adapted-process \dots
```

 $\label{eq:locale} \textit{nat-finite-adapted-process} \ \textit{M} \ \textit{F} \ \textit{0} \ :: \ \textit{nat} \ \textit{X} \ \textbf{for} \ \textit{M} \ \textit{F} \ \textit{X}$ 

 $\mbox{\bf locale}\ real\mbox{-} finite\mbox{-} adapted\mbox{-} process\ M\ F\ 0\ ::\ real\ X\ \mbox{\bf for}\ M\ F\ X$ 

sublocale nat-finite-adapted-process  $\subseteq$  nat-sigma-finite-adapted-process .. sublocale real-finite-adapted-process  $\subseteq$  real-sigma-finite-adapted-process ...

**locale** sigma-finite-adapted-process-order = sigma-finite-adapted-process  $M F t_0 X$  **for**  $M F t_0$  **and**  $X :: - \Rightarrow - \Rightarrow - :: \{order-topology, ordered-real-vector\}$ 

 $\label{locale} \textbf{locale} \ \ nat\text{-}sigma\text{-}finite\text{-}adapted\text{-}process\text{-}order = sigma\text{-}finite\text{-}adapted\text{-}process\text{-}order \\ M \ F \ 0 :: nat \ X \ \textbf{for} \ M \ F \ X$ 

 $\begin{tabular}{ll} \bf sublocale \ \it nat\mbox{-}\it sigma\mbox{-}\it finite\mbox{-}\it adapted\mbox{-}\it process\mbox{-}\it order \subseteq \it nat\mbox{-}\it sigma\mbox{-}\it finite\mbox{-}\it adapted\mbox{-}\it process\mbox{-}\it order \subseteq \it real\mbox{-}\it sigma\mbox{-}\it order \subseteq \it real\mbox{-}\it order \subseteq$ 

**locale** finite-adapted-process-order = finite-adapted-process  $M F t_0 X$  for  $M F t_0$  and  $X :: - \Rightarrow - \Rightarrow - :: \{order-topology, ordered-real-vector\}$ 

 $\label{locale} \begin{subarray}{l} \textbf{locale} \ nat\text{-}finite\text{-}adapted\text{-}process\text{-}order = finite\text{-}adapted\text{-}process\text{-}order M F 0 :: nat } X \ \textbf{for} \ M F X \end{subarray}$ 

 $\label{locale} \mbox{\bf locale real-finite-adapted-process-order} = \mbox{finite-adapted-process-order} \ M \ F \ 0 :: real \ X \ \mbox{\bf for} \ M \ F \ X$ 

 $\begin{tabular}{ll} \bf sublocale \ \it nat-finite-\it adapted-\it process-\it order \subseteq \it nat-\it sigma-\it finite-\it adapted-\it process-\it order \subseteq \it real-\it sigma-\it finite-\it adapted-\it process-\it order \subseteq \it real-\it sigma-\it finite-\it adapted-\it process-\it order \subseteq \it orde$ 

**locale** sigma-finite-adapted-process-linorder = sigma-finite-adapted-process-order  $M \ F \ t_0 \ X \ \text{for} \ M \ F \ t_0 \ \text{and} \ X :: - \Rightarrow - \Rightarrow - :: \{linorder-topology\}$ 

 $\begin{tabular}{l} \textbf{locale} \ nat\text{-}sigma\text{-}finite\text{-}adapted\text{-}process\text{-}linorder = sigma\text{-}finite\text{-}adapted\text{-}process\text{-}linorder \\ M\ F\ 0\ ::\ nat\ X\ \textbf{for}\ M\ F\ X \end{tabular}$ 

 $\label{locale} \textbf{locale} \ \textit{real-sigma-finite-adapted-process-linorder} = \textit{sigma-finite-adapted-process-linorder} \\ \textit{M} \ \textit{F} \ \textit{0} \ :: \ \textit{real} \ \textit{X} \ \textbf{for} \ \textit{M} \ \textit{F} \ \textit{X}$ 

 $\textbf{sublocale} \ \ nat\text{-}sigma\text{-}finite\text{-}adapted\text{-}process\text{-}linorder \subseteq nat\text{-}sigma\text{-}finite\text{-}adapted\text{-}process\text{-}order$  ..

```
\mathbf{sublocale}\ real\text{-}sigma\text{-}finite\text{-}adapted\text{-}process\text{-}linorder \subseteq real\text{-}sigma\text{-}finite\text{-}adapted\text{-}process\text{-}order
locale finite-adapted-process-linorder = finite-adapted-process-order M F t_0 X for
M F t_0 and X :: - \Rightarrow - \Rightarrow - :: \{linorder-topology\}
locale nat-finite-adapted-process-linorder = finite-adapted-process-linorder M F 0
:: nat \ X \ \mathbf{for} \ M \ F \ X
{\bf locale}\ real\mbox{-} finite\mbox{-} adapted\mbox{-} process\mbox{-} linorder\ =\ finite\mbox{-} adapted\mbox{-} process\mbox{-} linorder\ M\ F\ 0
:: real \ X \ \mathbf{for} \ M \ F \ X
\textbf{sublocale} \ \ nat\text{-}finite\text{-}adapted\text{-}process\text{-}linorder \subseteq nat\text{-}sigma\text{-}finite\text{-}adapted\text{-}process\text{-}linorder
\textbf{sublocale} \ real\text{-}finite\text{-}adapted\text{-}process\text{-}linorder \subseteq real\text{-}sigma\text{-}finite\text{-}adapted\text{-}process\text{-}linorder
7.2
               Martingale
locale martingale = sigma-finite-adapted-process +
   assumes integrable: \bigwedge i. t_0 \leq i \Longrightarrow integrable \ M \ (X \ i)
            and martingale-property: \bigwedge i \ j. \ t_0 \leq i \implies i \leq j \implies AE \ \xi \ in \ M. \ X \ i \ \xi =
cond-exp M (F i) (X j) \xi
locale martingale-order = martingale M F t_0 X for M F t_0 and X :: - \Rightarrow - \Rightarrow -
:: \{order-topology, ordered-real-vector\}
locale martingale-linorder = martingale M F t_0 X for M F t_0 and X :: - \Rightarrow - \Rightarrow
- :: {linorder-topology, ordered-real-vector}
sublocale martingale-linorder \subseteq martingale-order ..
lemma (in sigma-finite-filtered-measure) martingale-const-fun[intro]:
   assumes integrable M f f \in borel-measurable (F t_0)
   shows martingale M F t_0 (\lambda-. f)
  using assms sigma-finite-subalgebra.cond-exp-F-meas[OF - assms(1), THEN AE-symmetric]
borel-measurable-mono
   by (unfold-locales) blast+
lemma (in sigma-finite-filtered-measure) martingale-cond-exp[intro]:
   assumes integrable M f
   shows martingale M F t_0 (\lambda i. cond\text{-}exp M (F i) f)
   {\bf using} \ sigma-finite-subalgebra. borel-measurable-cond-exp' \ borel-
   by (unfold-locales) (auto intro: sigma-finite-subalgebra.cond-exp-nested-subalg[OF]
- assms] simp add: subalgebra-F subalgebras)
corollary (in sigma-finite-filtered-measure) martingale-zero[intro]: martingale\ M
F t_0 (\lambda - - \cdot \cdot \theta) by fastforce
corollary (in finite-filtered-measure) martingale-const[intro]: martingale M F t_0
```

 $(\lambda$ - -. c) by fastforce

### 7.3 Submartingale

```
locale \ submartingale = sigma-finite-adapted-process-order +
  assumes integrable: \bigwedge i. t_0 \leq i \Longrightarrow integrable \ M(Xi)
     cond\text{-}exp\ M\ (F\ i)\ (X\ j)\ \xi
locale submartingale-linorder = submartingale M F t_0 X for M F t_0 and X :: -
\Rightarrow - \Rightarrow - :: \{linorder-topology\}
sublocale martingale-order \subseteq submartingale using martingale-property by (unfold-locales)
(force simp add: integrable)+
sublocale martingale-linorder \subseteq submartingale-linorder ...
7.4
       Supermartingale
locale \ supermartingale = sigma-finite-adapted-process-order +
 assumes integrable: \bigwedge i. t_0 \leq i \Longrightarrow integrable \ M \ (X \ i)
     and supermartingale-property: \bigwedge i \ j. \ t_0 \leq i \Longrightarrow i \leq j \Longrightarrow AE \ \xi \ in \ M. \ X \ i \ \xi
\geq cond\text{-}exp\ M\ (F\ i)\ (X\ j)\ \xi
locale supermartingale-linorder = supermartingale M F t_0 X for M F t_0 and X
:: - \Rightarrow - \Rightarrow - :: \{linorder-topology\}
sublocale martingale-order \subseteq supermartingale using martingale-property by (unfold-locales)
(force simp add: integrable)+
sublocale martingale-linorder \subseteq supermartingale-linorder ...
A stochastic process is a martingale, if and only if it is both a submartingale
and a supermartingale.
lemma martingale-iff:
 shows martingale M F t_0 X \longleftrightarrow submartingale M F t_0 X \land supermartingale M
F t_0 X
proof (rule iffI)
 assume asm: martingale\ M\ F\ t_0\ X
 interpret martingale-order M F t_0 X by (intro martingale-order.intro asm)
  show submartingale M F t_0 X \wedge supermartingale M F t_0 X using submartin-
gale-axioms supermartingale-axioms by blast
next
  assume asm: submartingale M F t_0 X \wedge supermartingale M F t_0 X
 interpret submartingale M F t_0 X by (simp add: asm)
 interpret supermartingale M F t_0 X by (simp \ add: \ asm)
 show martingale M F t_0 X using submartingale-property supermartingale-property
```

#### 7.5 Martingale Lemmas

qed

**by** (unfold-locales) (intro integrable, blast, force)

In the following segment, we cover basic properties of martingales.

```
context martingale
begin
lemma cond-exp-diff-eq-zero:
 assumes t_0 \leq i \ i \leq j
 shows AE \xi in M. cond-exp M (F i) (\lambda \xi. X j \xi - X i \xi) \xi = 0
 using martingale-property[OF assms] assms
       sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i]
         sigma-finite-subalgebra.cond-exp-diff[OF-integrable(1,1), of Fiji] by
fast force
lemma set-integral-eq:
 assumes A \in F i t_0 \le i i \le j
 shows set-lebesgue-integral M A (X i) = set-lebesgue-integral M A (X j)
proof -
 interpret sigma-finite-subalgebra M F i using assms(2) by blast
  have \int x \in A. X \ i \ x \ \partial M = \int x \in A. cond-exp M (F \ i) (X \ j) x \ \partial M using
martingale-property[OF assms(2,3)] borel-measurable-cond-exp' assms subalgebras
subalgebra-def by (intro\ set-lebesgue-integral-cong-AE[OF - random-variable]) fast-
 also have ... = \int x \in A. X j x \partial M using assms by (auto simp: integrable intro:
cond-exp-set-integral[symmetric])
  finally show ?thesis.
qed
lemma scaleR-const[intro]:
 shows martingale M F t_0 (\lambda i \ x. \ c *_R X i \ x)
proof -
   fix i j :: 'b assume asm: t_0 \le i \ i \le j
   interpret sigma-finite-subalgebra M F i using asm by blast
    have AE \times in M. c *_R \times i \times x = cond\text{-}exp M (F i) (\lambda x. c *_R \times x j \times x) \times us
ing asm cond-exp-scaleR-right[OF integrable, of j, THEN AE-symmetric] martin-
gale-property[OF asm] by force
  }
 thus ?thesis by (unfold-locales) (auto simp add: integrable martingale.integrable)
qed
lemma uminus[intro]:
 shows martingale M F t_0 (-X)
 using scaleR-const[of -1] by (force\ intro:\ back-subst[of\ martingale\ M\ F\ t_0])
lemma add[intro]:
 assumes martingale M F t_0 Y
 shows martingale M F t_0 (\lambda i \xi. X i \xi + Y i \xi)
proof -
 interpret Y: martingale M F t_0 Y by (rule assms)
   fix i j :: 'b assume asm: t_0 \le i \ i \le j
```

```
hence AE \xi in M. X i \xi + Y i \xi = cond\text{-}exp M (F i) (<math>\lambda x. X j x + Y j x) \xi
       {\bf using} \ sigma-finite-subalgebra. cond-exp-add [OF-integrable\ martingale.integrable] OF-integrable\ martingale. integrable\ martingale. The property of 
assms], of F i j j, THEN AE-symmetric]
                       martingale-property[OF asm] martingale-martingale-property[OF assms
asm] by force
   thus ?thesis using assms
   by (unfold-locales) (auto simp add: integrable martingale.integrable)
qed
lemma diff[intro]:
   assumes martingale M F t_0 Y
   shows martingale M F t_0 (\lambda i x. X i x - Y i x)
proof -
   interpret Y: martingale M F t_0 Y by (rule assms)
       fix i j :: 'b assume asm: t_0 \leq i \ i \leq j
       hence AE \xi in M. X i \xi - Y i \xi = cond\text{-}exp M (F i) (<math>\lambda x. X j x - Y j x) \xi
       {f using}\ sigma-finite-subalgebra.cond-exp-diff[OF-integrable\ martingale.integrable[OF-integrable]]
assms], of F i j j, THEN AE-symmetric]
                       martingale-property[OF asm] martingale-martingale-property[OF assms
asm] by fastforce
   thus ?thesis using assms by (unfold-locales) (auto simp add: integrable martin-
gale.integrable)
qed
end
Using properties of the conditional expectation, we present the following
alternative characterizations of martingales.
lemma (in sigma-finite-adapted-process) martingale-of-cond-exp-diff-eq-zero:
   assumes integrable: \bigwedge i. t_0 \leq i \Longrightarrow integrable \ M \ (X \ i)
          and diff-zero: \bigwedge i \ j. t_0 \le i \Longrightarrow i \le j \Longrightarrow AE \ x \ in \ M. cond-exp M (F \ i) (\lambda \xi).
X j \xi - X i \xi) x = 0
       shows martingale M F t_0 X
proof
       fix i j :: 'b assume asm: t_0 \le i \ i \le j
       thus AE \xi in M. X i \xi = cond\text{-}exp M (F i) (X j) \xi
       using diff-zero [OF asm] sigma-finite-subalgebra. cond-exp-diff [OF - integrable (1,1),
of F i j i
                    sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] by
fastforce
qed (intro integrable)
lemma (in sigma-finite-adapted-process) martingale-of-set-integral-eq:
   assumes integrable: \bigwedge i. t_0 \leq i \Longrightarrow integrable \ M \ (X \ i)
```

```
and \bigwedge A \ i \ j. \ t_0 \leq i \Longrightarrow i \leq j \Longrightarrow A \in F \ i \Longrightarrow set-lebesgue-integral \ M \ A \ (X
i) = set-lebesgue-integral M A (X j)
      shows martingale M F t_0 X
proof (unfold-locales)
   fix i j :: 'b assume asm: t_0 \le i \ i \le j
   interpret sigma-finite-subalgebra M F i using asm by blast
  interpret r: sigma-finite-measure restr-to-subalq M (F i) by (simp add: sigma-fin-subalq)
      fix A assume A \in restr-to-subalg M (F i)
      hence *: A \in F i using sets-restr-to-subalg subalgebras asm by blast
    have set-lebesgue-integral (restr-to-subalg M(Fi)) A(Xi) = set-lebesgue-integral
MA(Xi) using * subalg asm by (auto simp: set-lebesque-integral-def intro: inte-
gral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator)
       also have ... = set-lebesgue-integral M A (cond-exp M (F i) (X j)) using *
assms(2)[OF asm] cond-exp-set-integral[OF integrable] asm by auto
    finally have set-lebesque-integral (restr-to-subaly M(F_i)) A(X_i) = set-lebesque-integral
(restr-to-subalq\ M\ (F\ i))\ A\ (cond-exp\ M\ (F\ i)\ (X\ j))\ using *subalq\ by\ (auto\ simp:
set-lebesque-integral-def intro!: integral-subalgebra2 [symmetric] borel-measurable-scaleR
borel-measurable-cond-exp borel-measurable-indicator)
   hence AE \ \xi in restr-to-subalq M \ (F \ i). X \ i \ \xi = cond-exp \ M \ (F \ i) \ (X \ j) \ \xi us-
ing asm by (intro r.density-unique-banach, auto intro: integrable-in-subalg subalg
borel-measurable-cond-exp integrable)
   thus AE \xi in M. X i \xi = cond\text{-}exp M (F i) (X j) \xi using AE\text{-}restr\text{-}to\text{-}subalg[OF]
subalg] by blast
qed (simp add: integrable)
7.6
             Submartingale Lemmas
context submartingale
begin
lemma cond-exp-diff-nonneg:
   assumes t_0 \leq i \ i \leq j
   shows AE x in M. cond-exp M (F i) (\lambda \xi. X j \xi - X i \xi) x \ge 0
  {f using} \ submarting ale	ext{-property} [OF\ assms] \ assms\ sigma-finite	ext{-subalgebra}. cond-exp-diff} [OF\ assms] 
- integrable(1,1), of - ji sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable
adapted, of i by fastforce
lemma add[intro]:
   assumes submartingale M F t_0 Y
   shows submartingale M F t_0 (\lambda i \xi. X i \xi + Y i \xi)
proof -
   interpret Y: submartingale\ M\ F\ t_0\ Y\ \mathbf{by}\ (rule\ assms)
      fix i j :: 'b assume asm: t_0 \leq i \ i \leq j
      hence AE \xi in M. X i \xi + Y i \xi \leq cond\text{-}exp M (F i) (<math>\lambda x. X j x + Y j x) \xi
      \mathbf{using}\ sigma-finite\text{-}subalgebra.cond\text{-}exp\text{-}add \lceil OF\text{-}integrable\ submartingale.integrable } \lceil OF\text{-}integrable\ submartingale.integrable\ submartingale.integrabl
assms], of F i j j]
```

```
submartingale-property[OF asm] submartingale.submartingale-property[OF
assms asm] add-mono[of X i - - Y i -] by force
   thus ?thesis using assms by (unfold-locales) (auto simp add: borel-measurable-add
random-variable adapted integrable Y.random-variable Y.adapted submartingale.integrable)
qed
lemma diff[intro]:
    assumes supermartingale M F t_0 Y
    shows submartingale M F t_0 (\lambda i \xi. X i \xi - Y i \xi)
    interpret Y: supermartingale M F t_0 Y by (rule assms)
        fix i j :: 'b assume asm: t_0 \le i \ i \le j
        hence AE \xi in M. X i \xi - Y i \xi \leq cond\text{-}exp M (F i) (<math>\lambda x. X j x - Y j x) \xi
         \textbf{using } \textit{sigma-finite-subalgebra}. cond-exp-diff [\textit{OF-integrable } \textit{supermartingale.integrable}] \textit{OF-integrable } \textit{supermartingale.integrable} [\textit{OF-integrable } \textit{supermartingale.integrable}] \textit{OF-integrable } \textit{supermartingale.integrable} \textit{oF-integrable} \textit{oF-integr
assms], of F i j j]
                  submartingale-property[OF asm] supermartingale-supermartingale-property[OF]
assms asm diff-mono of X i - - Y i by force
   thus ?thesis using assms by (unfold-locales) (auto simp add: borel-measurable-diff
random-variable adapted integrable Y.random-variable Y.adapted supermartingale.integrable)
qed
lemma scaleR-nonneg:
    assumes c \geq \theta
   shows submartingale M F t_0 (\lambda i \xi. c *_R X i \xi)
proof
        fix i j :: 'b assume asm: t_0 \le i \ i \le j
        thus AE \xi in M. c *_R X i \xi \leq cond\text{-}exp M (F i) (\lambda \xi. c *_R X j \xi) \xi
                \mathbf{using} \ \mathit{sigma-finite-subalgebra}. \mathit{cond-exp-scaleR-right}[\mathit{OF-integrable}, \ \mathit{of} \ \mathit{F} \ \mathit{i}
j c] submartingale-property[OF asm] by (fastforce intro!: scaleR-left-mono[OF -
assms)
    }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)
\mathbf{lemma}\ scaleR-le-zero:
    assumes c \leq \theta
    shows supermartingale M F t_0 (\lambda i \xi. c *_R X i \xi)
proof
    {
        fix i j :: 'b assume asm: t_0 \le i \ i \le j
        thus AE \xi in M. c *_R X i \xi \geq cond\text{-}exp M (F i) (\lambda \xi. c *_R X j \xi) \xi
              using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of F i j
c] submartingale-property[OF asm]
```

```
by (fastforce intro!: scaleR-left-mono-neg[OF - assms])
 }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)
lemma uminus[intro]:
  shows supermartingale M F t_0 (-X)
  unfolding fun-Compl-def using scaleR-le-zero[of -1] by simp
end
context submartingale-linorder
begin
lemma set-integral-le:
 assumes A \in F \ i \ t_0 \le i \ i \le j
 shows set-lebesgue-integral M A (X i) \leq set-lebesgue-integral M A (X j)
 using submartingale-property[OF assms(2), of j] assms subalgebras
  by (subst sigma-finite-subalgebra.cond-exp-set-integral [OF - integrable \ assms(1),
of j
   (auto\ intro!:\ scaleR-left-mono\ integral-mono-AE-banach\ integrable-mult-indicator
integrable simp add: subalgebra-def set-lebesgue-integral-def)
lemma max:
 assumes submartingale-linorder M F t_0 Y
 shows submartingale-linorder M F t_0 (\lambda i \xi. max (X i \xi) (Y i \xi))
proof (unfold-locales)
 interpret Y: submartingale-linorder M F t<sub>0</sub> Y by (rule assms)
   fix i j :: 'b assume asm: t_0 \leq i \ i \leq j
    have AE \xi in M. max (X i \xi) (Y i \xi) \leq max (cond-exp M (F i) (X j) \xi)
(cond\text{-}exp\ M\ (F\ i)\ (Y\ j)\ \xi) using submartingale\text{-}property\ Y.submartingale\text{-}property
asm unfolding max-def by fastforce
   thus AE \xi in M. max (X i \xi) (Y i \xi) \leq cond\text{-}exp M (F i) (\lambda \xi. max (X j \xi) (Y i \xi))
j \in \{1\} using sigma-finite-subalgebra.cond-exp-max OF - integrable Y.integrable, of
F \ i \ j \ j asm by (fast intro: order.trans)
  show \bigwedge i. t_0 \leq i \Longrightarrow (\lambda \xi. \ max \ (X \ i \ \xi) \ (Y \ i \ \xi)) \in borel-measurable \ (F \ i) \ \bigwedge i.
t_0 \leq i \implies integrable \ M \ (\lambda \xi. \ max \ (X \ i \ \xi) \ (Y \ i \ \xi)) \ \mathbf{by} \ (force \ intro: \ Y.integrable
integrable \ assms)+
qed
lemma max-\theta:
 shows submartingale-linorder M F t_0 (\lambda i \xi. max \theta (X i \xi))
proof -
  interpret zero: martingale-linorder M F t_0 \lambda- -. 0 by (force intro: martin-
gale-linorder.intro martingale-order.intro)
 show ?thesis by (intro zero.max submartingale-linorder.intro submartingale-axioms)
qed
```

#### end

```
lemma (in sigma-finite-adapted-process-order) submartingale-of-cond-exp-diff-nonneg:
  assumes integrable: \bigwedge i. t_0 \leq i \implies integrable M(X i)
      and diff-nonneg: \bigwedge i \ j. t_0 \le i \implies i \le j \implies AE \ x \ in \ M. cond-exp M (F \ i)
(\lambda \xi. \ X \ j \ \xi - X \ i \ \xi) \ x \ge 0
   shows submartingale M F t_0 X
proof (unfold-locales)
   fix i j :: 'b assume asm: t_0 \le i i \le j
   thus AE \xi in M. X i \xi \leq cond\text{-}exp M (F i) (X j) \xi
      using diff-nonneg[OF asm] sigma-finite-subalgebra.cond-exp-diff[OF - inte-
grable(1,1), of F i j i
          sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] by
fast force
qed (intro integrable)
lemma (in sigma-finite-adapted-process-linorder) submartingale-of-set-integral-le:
 assumes integrable: \bigwedge i. t_0 \leq i \Longrightarrow integrable \ M \ (X \ i)
     and \bigwedge A \ i \ j. \ t_0 \le i \Longrightarrow i \le j \Longrightarrow A \in F \ i \Longrightarrow set-lebesgue-integral \ M \ A \ (X
i) \leq set-lebesgue-integral M \land (X \ j)
   shows submartingale M F t_0 X
proof (unfold-locales)
   fix i j :: 'b assume asm: t_0 \le i i \le j
  interpret r: sigma-finite-measure restr-to-subalq M (Fi) using asm sigma-finite-subalqebra.sigma-fin-subalq
\mathbf{by} blast
    {
     fix A assume A \in restr-to-subalg M (F i)
     hence *: A \in F i using asm sets-restr-to-subalg subalgebras by blast
    have set-lebesgue-integral (restr-to-subalg M(Fi)) A(Xi) = set-lebesgue-integral
M A (X i) using * asm subalgebras by (auto simp: set-lebesgue-integral-def intro:
integral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator)
      also have ... \leq set-lebesgue-integral M A (cond-exp M (F i) (X j)) using
* assms(2)[OF \ asm] \ asm \ sigma-finite-subalgebra.cond-exp-set-integral[OF - inte-subalgebra]
grable] by fastforce
     also have ... = set-lebesque-integral (restr-to-subalg M (F i)) A (cond-exp M
(F \ i) \ (X \ j)) using * asm subalgebras by (auto simp: set-lebesque-integral-def intro!:
integral-subalgebra2[symmetric] \ borel-measurable-scaleR \ borel-measurable-cond-exp
borel-measurable-indicator)
    finally have 0 \le set-lebesque-integral (restr-to-subalg M (F i)) A (\lambda \xi. cond-exp
M(F i)(X j) \xi - X i \xi) using * asm subalgebras by (subst set-integral-diff,
auto simp add: set-integrable-def sets-restr-to-subalg intro!: integrable adapted inte-
grable\mbox{-}in\mbox{-}subalg\ borel\mbox{-}measurable\mbox{-}scaleR\ borel\mbox{-}measurable\mbox{-}indicator\ borel\mbox{-}measurable\mbox{-}cond\mbox{-}exp
integrable-mult-indicator)
   hence AE \xi in restr-to-subalg M(Fi). 0 \le cond\text{-}exp\ M(Fi)(Xj)\xi - Xi\xi
```

```
by (intro r.density-nonneq integrable-in-subalq asm subalqebras borel-measurable-diff
borel-measurable-cond-exp\ adapted\ Bochner-Integration.integrable-diff\ integrable-cond-exp
integrable)
           thus AE \xi in M. Xi \xi \leq cond\text{-}exp M (Fi) (Xj) \xi using AE\text{-}restr\text{-}to\text{-}subalg[OF]
subalgebras asm by simp
qed (intro integrable)
                                 Supermartingale Lemmas
7.7
The following lemmas are exact duals of the ones for submartingales.
context supermartingale
begin
lemma cond-exp-diff-nonneg:
       assumes t_0 \leq i \ i \leq j
       shows AE x in M. cond-exp M (F i) (\lambda \xi. X i \xi - X j \xi) x \ge 0
     \textbf{using} \ assms \ supermarting ale-property [OF \ assms] \ sigma-finite-subalgebra. cond-exp-diff [OF \ assms] \ sigma-finite-subalgebra. \ sigma-finite-
 - integrable(1,1), of F i i j
                                    sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] by
fast force
lemma add[intro]:
       assumes supermartingale M F t_0 Y
       shows supermartingale M F t_0 (\lambda i \xi. X i \xi + Y i \xi)
        interpret Y: supermartingale M F t_0 Y by (rule assms)
               fix i j :: 'b assume asm: t_0 \le i \ i \le j
               hence AE \xi in M. X i \xi + Y i \xi \geq cond\text{-}exp M (F i) (<math>\lambda x. X j x + Y j x) \xi
                \textbf{using } \textit{sigma-finite-subalgebra}. cond-exp-add [\textit{OF-integrable } \textit{supermartingale.integrable}] \textit{OF-integrable } \textit{Supermartingale.integrable} \textit{Sup
assms, of F i j j
                                 supermarting a le-property [OF\ asm]\ supermarting a le-supermarting a le-property [OF\ asm]\ supermarting a le-property [OF\ asm]\ supermarting
assms asm] add-mono[of - Xi - - Yi -] by force
    thus ?thesis using assms by (unfold-locales) (auto simp add: borel-measurable-add
random-variable adapted integrable Y-random-variable Y-adapted supermartingale.integrable)
qed
lemma diff[intro]:
       assumes submartingale M F t_0 Y
       shows supermartingale M F t_0 (\lambda i \xi. X i \xi - Y i \xi)
proof -
       interpret Y: submartingale\ M\ F\ t_0\ Y\ \mathbf{by}\ (rule\ assms)
               fix i j :: 'b assume asm: t_0 \leq i \ i \leq j
               hence AE \xi in M. X i \xi - Y i \xi \ge cond\text{-}exp M (F i) (<math>\lambda x. X j x - Y j x) \xi
```

 $\textbf{using } \textit{sigma-finite-subalgebra}. \textit{cond-exp-diff} \, [\textit{OF-integrable submartingale.integrable} \, [\textit{OF-integrable submartingale.integrable}] \\$ 

```
assms], of F i j j, unfolded fun-diff-def]
       supermarting a le-property [OF\ asm]\ submarting a le. submarting a le-property [OF\ asm]
assms asm] diff-mono[of - X i - Y i -] by force
 thus ?thesis using assms by (unfold-locales) (auto simp add: borel-measurable-diff
random-variable adapted integrable Y.random-variable Y.adapted submartingale.integrable)
qed
lemma scaleR-nonneg:
 assumes c \geq \theta
 shows supermartingale M F t_0 (\lambda i \, \xi. \, c *_R X \, i \, \xi)
proof
   fix i j :: 'b assume asm: t_0 \le i \ i \le j
   thus AE \xi in M. c *_R X i \xi \geq cond\text{-}exp M (F i) (<math>\lambda \xi. c *_R X j \xi) \xi
      using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of F i
j c] supermartingale-property[OF asm] by (fastforce intro!: scaleR-left-mono[OF -
assms])
 }
{\bf qed}\ (auto\ simp\ add:\ borel-measurable-integrable\ borel-measurable-scaleR\ integrable
random-variable adapted borel-measurable-const-scaleR)
\mathbf{lemma} scaleR-le-zero:
 assumes c \leq \theta
 shows submartingale M F t_0 (\lambda i \xi. c *_R X i \xi)
proof
  {
   fix i j :: 'b assume asm: t_0 \leq i \ i \leq j
   thus AE \xi in M. c *_R X i \xi \leq cond\text{-}exp M (F i) (<math>\lambda \xi. c *_R X j \xi) \xi
    using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of F i j c]
supermarting ale-property[OF\ asm]\ by (fast force\ intro!:\ scaleR-left-mono-neg[OF\ -
assms])
 }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)
lemma uminus[intro]:
 shows submartingale M F t_0 (-X)
 unfolding fun-Compl-def using scaleR-le-zero[of -1] by simp
end
{f context} supermarting a le-linor der
begin
lemma set-integral-ge:
 assumes A \in F i t_0 \le i i \le j
 shows set-lebesgue-integral M A (X i) \geq set-lebesgue-integral M A (X j)
```

```
by (subst\ sigma-finite-subalgebra.\ cond-exp-set-integral[OF-integrable\ assms(1),
of j])
   (auto\ intro!:\ scaleR-left-mono\ integral-mono-AE-banach\ integrable-mult-indicator
integrable simp add: subalgebra-def set-lebesque-integral-def)
lemma min:
  assumes supermartingale-linorder M F t_0 Y
  shows supermartingale-linorder M F t_0 (\lambda i \xi. min (X i \xi) (Y i \xi))
proof (unfold-locales)
  interpret Y: supermartingale-linorder M F t_0 Y by (rule assms)
   fix i j :: 'b assume asm: t_0 \leq i \ i \leq j
  have AE \xi in M. min(X i \xi)(Y i \xi) \ge min(cond-exp M(F i)(X j)\xi)(cond-exp
M\ (F\ i)\ (Y\ j)\ \xi) using supermartingale-property Y.supermartingale-property asm
unfolding min-def by fastforce
   thus AE \xi in M. min (X i \xi) (Y i \xi) \ge cond\text{-}exp M (F i) (\lambda \xi. min (X j \xi) (Y i \xi))
j \in \{1, 1\} \in \{1, 1\} using sigma-finite-subalgebra.cond-exp-min OF - integrable Y.integrable, of
F \ i \ j \ j asm by (fast intro: order.trans)
  show \bigwedge i. \ t_0 \leq i \Longrightarrow (\lambda \xi. \ min \ (X \ i \ \xi) \ (Y \ i \ \xi)) \in borel-measurable \ (F \ i) \ \bigwedge i.
t_0 \leq i \implies integrable \ M \ (\lambda \xi. \ min \ (X \ i \ \xi) \ (Y \ i \ \xi)) \ \mathbf{by} \ (force \ intro: \ Y.integrable
integrable \ assms)+
qed
lemma min-\theta:
 shows supermartingale-linorder M F t_0 (\lambda i \ \xi. min \theta (X i \ \xi))
proof -
  interpret zero: martingale-linorder M F t_0 \lambda- -. \theta by (force intro: martin-
gale-linorder.intro)
  show ?thesis by (intro zero.min supermartingale-linorder.intro supermartin-
gale-axioms)
qed
end
lemma (in sigma-finite-adapted-process-order) supermartingale-of-cond-exp-diff-le-zero:
  assumes integrable: \bigwedge i. t_0 \leq i \Longrightarrow integrable \ M \ (X \ i)
      and diff-le-zero: \bigwedge i \ j. \ t_0 \le i \Longrightarrow i \le j \Longrightarrow AE \ x \ in \ M. \ cond-exp \ M \ (F \ i)
(\lambda \xi. \ X \ j \ \xi - X \ i \ \xi) \ x \le 0
   shows supermartingale M F t_0 X
proof
  {
   fix i j :: 'b assume asm: t_0 \le i \ i \le j
   thus AE \xi in M. X i \xi \geq cond\text{-}exp M (F i) (X j) \xi
       using diff-le-zero[OF asm] sigma-finite-subalgebra.cond-exp-diff[OF - inte-
grable(1,1), of F i j i
           sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] by
fast force
```

using supermartingale-property[OF assms(2), of j]  $assms\ subalgebras$ 

```
qed (intro integrable)
lemma (in sigma-finite-adapted-process-linorder) supermartingale-of-set-integral-ge:
 assumes integrable: \bigwedge i. t_0 \leq i \Longrightarrow integrable \ M \ (X \ i)
     and \bigwedge A \ i \ j. \ t_0 \leq i \Longrightarrow i \leq j \Longrightarrow A \in F \ i \Longrightarrow set-lebesgue-integral \ M \ A \ (X
j) \leq set-lebesgue-integral M \land (X \mid i)
   shows supermartingale M F t_0 X
proof -
 interpret -: adapted-process M F t_0 - X by (rule uminus-adapted)
 interpret uminus-X: sigma-finite-adapted-process-linorder M F t_0 -X ..
 \mathbf{note} * = set\text{-}integral\text{-}uminus[unfolded\ set\text{-}integrable\text{-}def,\ OF\ integrable\text{-}mult\text{-}indicator[OF\ ]}
- integrable]]
 have supermartingale M F t_0 (-(-X))
  \mathbf{using}\ ord\text{-}eq\text{-}le\text{-}trans[OF* ord\text{-}le\text{-}eq\text{-}trans[OF le\text{-}imp\text{-}neg\text{-}le[OF\ assms(2)]*}[symmetric]]]
subalgebras
   by (intro submartingale.uminus uminus-X.submartingale-of-set-integral-le)
      (clarsimp simp add: fun-Compl-def subalgebra-def integrable | fastforce)+
 thus ?thesis unfolding fun-Compl-def by simp
qed
Many of the statements we have made concerning martingales can be sim-
plified when the indexing set is the natural numbers. Given a point in time
i \in \mathbb{N}, it suffices to consider the successor i + (1::'a), instead of all future
times i \leq j.
7.8
       Discrete Time Martingales
locale nat-martingale = martingale M F 0 :: nat X for M F X
locale nat-submartingale = submartingale M F 0 :: nat X \text{ for } M F X
locale nat-supermartingale = supermartingale M F \theta :: nat X for M F X
locale nat-submartingale-linorder = submartingale-linorder M F 0 :: nat X for M
FX
locale nat-supermartingale-linorder = supermartingale-linorder M F 0 :: nat X
for M F X
sublocale nat-submartingale-linorder \subseteq nat-submartingale ...
sublocale nat-supermartingale-linorder \subseteq nat-supermartingale ...
A predictable martingale is necessarily constant.
lemma (in nat-martingale) predictable-const:
 assumes nat-predictable-process M F X
 shows AE \xi in M. X i \xi = X j \xi
 have *: AE \xi in M. X i \xi = X \theta \xi  for i
 proof (induction i)
   case \theta
   then show ?case by (simp add: bot-nat-def)
```

```
next
      case (Suc\ i)
    interpret S: nat-adapted-process MF \lambda i. X (Suc i) by (intro nat-predictable-process.adapted-Suc
       show ?case using Suc S.adapted[of i] martingale-property[OF - le-SucI, of i]
sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable, of F i Suc i] by fastforce
   show ?thesis using *[of i] *[of j] by force
qed
lemma (in nat-sigma-finite-adapted-process) martingale-of-set-integral-eq-Suc:
   assumes integrable: \bigwedge i. integrable M(X i)
         and \bigwedge A \ i. \ A \in F \ i \Longrightarrow set-lebesgue-integral M \ A \ (X \ i) = set-lebesgue-integral
M A (X (Suc i))
      shows nat-martingale M F X
proof (intro nat-martingale.intro martingale-of-set-integral-eq)
   fix i j A assume asm: i \leq j A \in sets (F i)
   show set-lebesgue-integral M A (X i) = set-lebesgue-integral M A (X j) using
   proof (induction j - i arbitrary: i j)
      case \theta
      then show ?case using asm by simp
   next
      case (Suc \ n)
      hence *: n = j - Suc \ i \ \mathbf{by} \ linarith
      have Suc\ i \leq j using Suc(2,3) by linarith
       thus ?case using sets-F-mono[OF - le-SucI] Suc(4) Suc(1)[OF *] by (auto
intro: assms(2)[THEN trans])
   qed
qed (simp add: integrable)
lemma (in nat-sigma-finite-adapted-process) martingale-nat:
   assumes integrable: \bigwedge i. integrable M(X i)
          and \bigwedge i. AE \xi in M. X i \xi = cond-exp M (F i) (X (Suc i)) \xi
      shows nat-martingale M F X
proof (unfold-locales)
   fix i j :: nat assume asm: i \leq j
   show AE \xi in M. X i \xi = cond\text{-}exp M (F i) (X j) \xi using asm
   proof (induction j - i arbitrary: i j)
      case \theta
      hence j = i by simp
     \textbf{thus}~? case~\textbf{using}~sigma-finite-subalgebra.cond-exp-F-meas [OF-integrable~adapted, and of the control of
 THEN AE-symmetric by blast
   next
      case (Suc \ n)
      have j: j = Suc (n + i) using Suc by linarith
      have n: n = n + i - i using Suc by linarith
      have *: AE \xi in M. cond\text{-}exp M (F (n + i)) (X j) \xi = X (n + i) \xi  unfolding
j using assms(2)[THEN\ AE-symmetric] by blast
```

```
have AE \xi in M. cond-exp M (F i) (X j) \xi = cond-exp M (F i) (cond-exp M)
(F(n+i))(Xj)) \xi by (intro cond-exp-nested-subalg integrable subalg, simp add:
subalgebra-def sets-F-mono)
      hence AE \xi in M. cond-exp M (F i) (X j) \xi = cond-exp M (F i) <math>(X (n + i))
\xi using cond-exp-cong-AE[OF integrable-cond-exp integrable *] by force
      thus ?case using Suc(1)[OF\ n] by fastforce
   qed
qed (simp add: integrable)
\mathbf{lemma} \ (\mathbf{in} \ nat\text{-}sigma\text{-}finite\text{-}adapted\text{-}process) \ martingale\text{-}of\text{-}cond\text{-}exp\text{-}diff\text{-}Suc\text{-}eq\text{-}zero:
   assumes integrable: \bigwedge i. integrable M(X i)
          and \bigwedge i. AE \xi in M. cond-exp M (F i) (\lambda \xi. X (Suc i) \xi – X i \xi) \xi = 0
      shows nat-martingale M F X
proof (intro martingale-nat integrable)
   \mathbf{fix} i
  show AE \ \xi \ in \ M. \ Xi \ \xi = cond-exp \ M \ (Fi) \ (X \ (Suc \ i)) \ \xi \ using \ cond-exp-diff[OF]
integrable(1,1), of i Suc i i | cond-exp-F-meas[OF integrable adapted, of i] assms(2)[of integrable(1,1), of i Suc i i | cond-exp-F-meas[OF integrable adapted, of i] assms(2)[of integrable(1,1), of i Suc i i | cond-exp-F-meas[OF integrable adapted, of i] assms(2)[of integrable(1,1), of i Suc i i | cond-exp-F-meas[OF integrable adapted, of i] assms(2)[of integrable adapt
i] by fastforce
qed
7.9
              Discrete Time Submartingales
lemma (in nat-submartingale) predictable-mono:
   assumes nat-predictable-process M F X i \leq j
   shows AE \xi in M. X i \xi \leq X j \xi
   using assms(2)
proof (induction j - i arbitrary: i j)
   case \theta
   then show ?case by simp
next
   case (Suc\ n)
   hence *: n = j - Suc \ i by linarith
  interpret S: nat-adapted-process M F \lambda i. X (Suc i) by (intro nat-predictable-process adapted-Suc
   have Suc\ i \leq j using Suc(2,3) by linarith
    thus ?case using Suc(1)[OF *] S.adapted[of i] submartingale-property[OF -]
le-SucI, of i] sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable, of F i Suc
i by fastforce
qed
lemma (in nat-sigma-finite-adapted-process-linorder) submartingale-of-set-integral-le-Suc:
   assumes integrable: \bigwedge i. integrable M(X i)
         and \bigwedge A \ i. \ A \in F \ i \Longrightarrow set-lebesgue-integral M \ A \ (X \ i) \le set-lebesgue-integral
M A (X (Suc i))
      shows nat-submartingale M F X
proof (intro nat-submartingale.intro submartingale-of-set-integral-le)
   fix i j A assume asm: i \leq j A \in sets (F i)
   show set-lebesgue-integral M A (X \ i) \leq set-lebesgue-integral M A (X \ j) using
```

```
proof (induction j - i arbitrary: i j)
         case \theta
         then show ?case using asm by simp
     next
         case (Suc\ n)
         hence *: n = j - Suc \ i \ \mathbf{by} \ linarith
         have Suc\ i \leq j using Suc(2,3) by linarith
          thus ?case using sets-F-mono[OF - le-SucI] Suc(4) Suc(1)[OF *] by (auto
intro: assms(2)[THEN \ order-trans])
     qed
qed (simp add: integrable)
lemma (in nat-sigma-finite-adapted-process-linorder) submartingale-nat:
     assumes integrable: \bigwedge i. integrable M(X i)
             and \bigwedge i. AE \xi in M. X i \xi < cond\text{-}exp\ M\ (F\ i)\ (X\ (Suc\ i))\ \xi
         shows nat-submartingale M F X
    using subalq integrable assms(2)
   \textbf{by} \ (intro\ submarting ale-of-set-integral-le-Suc\ ord-le-eq-trans[OF\ set-integral-mono-AE-banach]) and the submarting ale-of-set-integral-le-suc\ ord-le-suc\ ord-le-
cond-exp-set-integral[symmetric]], simp)
         (meson in-mono integrable-mult-indicator set-integrable-def subalgebra-def, me-
son integrable-cond-exp in-mono integrable-mult-indicator set-integrable-def subal-
gebra-def, fast+)
lemma (in nat-sigma-finite-adapted-process-linorder) submartingale-of-cond-exp-diff-Suc-nonneg:
     assumes integrable: \bigwedge i. integrable M(X i)
             and \bigwedge i. AE \xi in M. cond-exp M (F i) (\lambda \xi. X (Suc i) \xi – X i \xi) \xi \geq 0
         shows nat-submartingale M F X
proof (intro submartingale-nat integrable)
    \mathbf{fix} i
   show AE \xi in M. X i \xi \leq cond\text{-}exp M (F i) (X (Suc i)) \xi using cond\text{-}exp-diff[OF]
integrable(1,1), of i Suc i i] cond-exp-F-meas[OF integrable adapted, of i] assms(2)[of integrable adapted
i by fastforce
\mathbf{qed}
lemma (in nat-submartingale-linorder) partial-sum-scaleR:
     assumes nat-adapted-process M F C \wedge i. AE \xi in M. 0 < C i \xi \wedge i. AE \xi in
M. Ci \xi \leq R
    shows nat-submartingale M F (\lambda n \xi. \sum i < n. C i \xi *_R (X (Suc i) \xi - X i \xi))
proof-
    interpret C: nat-adapted-process M F C by (rule assms)
     interpret C': nat-adapted-process M F \lambda i \xi. C (i-1) \xi *_R (X i \xi - X (i \xi - X \xi))
(-1) \xi by (intro nat-adapted-process.intro adapted-process.scaleR-right-adapted
adapted-process. diff-adapted, unfold-locales) (auto intro: adaptedD C. adaptedD)+
    interpret C'': nat-adapted-process M F \lambda n \xi. \sum i < n. C i \xi *_R (X (Suc i) \xi - i)
X \ i \ \xi) by (rule C'.partial-sum-Suc-adapted[unfolded diff-Suc-1])
   interpret S: nat-sigma-finite-adapted-process-linorder M F (\lambda n \xi. \sum i < n. C i \xi
*_R (X (Suc i) \xi - X i \xi)) ...
  have integrable M (\lambda x. C i x *_R (X (Suc\ i)\ x - X\ i\ x)) for i using assms(2,3)[of
i] by (intro Bochner-Integration.integrable-bound[OF integrable-scaleR-right, OF
```

```
Bochner-Integration.integrable-diff, OF integrable (1,1), of R Suc i i) (auto simp
add: mult-mono)
 moreover have AE \xi in M. 0 \leq cond-exp M (F i) (\lambda \xi. (\sum i < Suc i. C i \xi *_R
(X (Suc \ i) \ \xi - X \ i \ \xi)) - (\sum i < i. \ C \ i \ \xi *_R (X (Suc \ i) \ \xi - X \ i \ \xi))) \ \xi  for i
     using sigma-finite-subalgebra.cond-exp-measurable-scale R[OF-calculation-
C.adapted, of i
         cond-exp-diff-nonneg[OF - le-SucI, OF - order.refl, of i] assms(2,3)[of\ i]
by (fastforce simp add: scaleR-nonneg-nonneg integrable)
 ultimately show ?thesis by (intro S.submartingale-of-cond-exp-diff-Suc-nonneg
Bochner-Integration.integrable-sum, blast+)
qed
lemma (in nat-submartingale-linorder) partial-sum-scaleR':
 assumes nat-predictable-process M F C \wedge i. AE \xi in M. 0 \leq C i \xi \wedge i. AE \xi in
M. Ci \xi \leq R
 shows nat-submartingale M F (\lambda n \xi. \sum i < n. C (Suc i) \xi *_R (X (Suc i) \xi - X)
i \xi)
proof
 interpret C: nat-predictable-process M F C by (rule assms)
 interpret Suc-C: nat-adapted-process M F \lambda i. C (Suc i) using C.adapted-Suc.
 show ?thesis by (intro partial-sum-scaleR[of - R] assms) (intro-locales)
qed
         Discrete Time Supermartingales
7.10
lemma (in nat-supermartingale) predictable-mono:
 assumes nat-predictable-process M F X i \leq j
 shows AE \xi in M. X i \xi \geq X j \xi
 using assms(2)
proof (induction j - i arbitrary: i j)
 case \theta
 then show ?case by simp
next
 case (Suc \ n)
 hence *: n = j - Suc \ i \ \mathbf{by} \ linarith
 interpret S: nat-adapted-process M F \lambda i. X (Suc i) by (intro nat-predictable-process.adapted-Suc
assms)
 have Suc \ i < i \ using \ Suc(2,3) by linarith
  thus ?case using Suc(1)[OF *] S.adapted[of i] supermartingale-property[OF -
le-SucI, of i | sigma-finite-subalgebra.cond-exp-F-meas|OF - integrable, of F i Suc
i by fastforce
qed
{\bf lemma\ (in\ }nat\text{-}sigma\text{-}finite\text{-}adapted\text{-}process\text{-}linorder)\ supermarting ale\text{-}of\text{-}set\text{-}integral\text{-}ge\text{-}Suc:}
 assumes integrable: \bigwedge i. integrable M(X i)
    and \bigwedge A \ i. \ A \in F \ i \Longrightarrow set-lebesgue-integral M \ A \ (X \ i) \ge set-lebesgue-integral
M A (X (Suc i))
   shows nat-supermartingale M F X
proof -
```

```
interpret -: adapted-process M F \theta - X by (rule uminus-adapted)
  interpret uminus-X: nat-sigma-finite-adapted-process-linorder M F - X..
 \mathbf{note} * = set\text{-}integral\text{-}uminus[unfolded\ set\text{-}integrable\text{-}def,\ OF\ integrable\text{-}mult\text{-}indicator[OF\ ]}
- integrable]]
  have nat-supermartingale M F (-(-X))
  \mathbf{using}\ ord\text{-}eq\text{-}le\text{-}trans[OF* ord\text{-}le\text{-}eq\text{-}trans[OF le\text{-}imp\text{-}neq\text{-}le[OF\ assms(2)]*}[symmetric]]]
subalgebras
  by (intro nat-supermartingale.intro submartingale.uminus nat-submartingale.axioms
uminus-X.submartingale-of-set-integral-le-Suc)
       (clarsimp\ simp\ add:\ fun\mbox{-}Compl\mbox{-}def\ subalgebra\mbox{-}def\ integrable\ |\ fastforce) +
  thus ?thesis unfolding fun-Compl-def by simp
qed
lemma (in nat-sigma-finite-adapted-process-linorder) supermartingale-nat:
  assumes integrable: \bigwedge i. integrable M(X i)
     and \bigwedge i. AE \xi in M. X i \xi > cond\text{-}exp\ M\ (F\ i)\ (X\ (Suc\ i))\ \xi
   shows nat-supermartingale M F X
proof -
  interpret -: adapted-process M F \theta - X by (rule uminus-adapted)
  interpret uminus-X: nat-sigma-finite-adapted-process-linorder M F - X...
 have AE \xi in M. -Xi \xi \leq cond\text{-}exp\ M\ (Fi)\ (\lambda x. -X\ (Suc\ i)\ x)\ \xi for i using
assms(2) cond-exp-uminus[OF integrable, of i Suc i] by force
  hence nat-supermartingale M F (-(-X)) by (intro nat-supermartingale.intro
submarting a le. uminus\ nat-submarting a le. axioms\ uminus-X. submarting a le-nat)\ (automarting a le. uminus\ nat-submarting a le. axioms\ uminus-X. submarting a le-nat)
simp add: fun-Compl-def integrable)
  thus ?thesis unfolding fun-Compl-def by simp
qed
{\bf lemma~(in~} nat\text{-}sigma\text{-}finite\text{-}adapted\text{-}process\text{-}linorder)~supermartingale\text{-}of\text{-}cond\text{-}exp\text{-}diff\text{-}Suc\text{-}le\text{-}zero:}
 assumes integrable: \bigwedge i. integrable M(X i)
     and \bigwedge i. AE \xi in M. cond-exp M (F i) (\lambda \xi. X (Suc i) \xi – X i \xi) \xi \leq 0
   shows nat-supermartingale M F X
proof (intro supermartingale-nat integrable)
 show AE \xi in M. X i \xi \ge cond\text{-}exp M (F i) (X (Suc i)) \xi using cond\text{-}exp-diff[OF]
integrable (1,1), of i Suc i i] cond-exp-F-meas[OF integrable adapted, of i] assms(2)[of
i] by fastforce
qed
end
theory Example-Coin-Toss
 imports Martingale HOL-Probability. Stream-Space HOL-Probability. Probability-Mass-Function
begin
```

## 8 Example: Coin Toss

We consider a probability space consisting of infinite sequences of coin tosses.

```
definition bernoulli-stream :: real \Rightarrow (bool stream) measure where bernoulli-stream p = stream-space (measure-pmf (bernoulli-pmf p))
```

**lemma** space-bernoulli-stream[simp]: space (bernoulli-stream p) = UNIV by (simp add: bernoulli-stream-def space-stream-space)

We define the fortune of the player at time n to be the number of heads minus number of tails.

```
definition fortune :: nat \Rightarrow bool \ stream \Rightarrow real \ \mathbf{where} fortune n = (\lambda s. \sum b \leftarrow stake \ (Suc \ n) \ s. \ if \ b \ then \ 1 \ else \ -1)
definition toss :: nat \Rightarrow bool \ stream \Rightarrow real \ \mathbf{where} toss \ n = (\lambda s. \ if \ snth \ s \ n \ then \ 1 \ else \ -1)
```

lemma toss-indicator-def: toss  $n = indicator \{s. s !! n\} - indicator \{s. \neg s !! n\}$  unfolding toss-def indicator-def by force

```
lemma range-toss: range (toss\ n) = \{-1,\ 1\}
proof –
have sconst True !! n by simp
moreover have \neg sconst\ False !! n by simp
ultimately have \exists\ x.\ x !! n\ \exists\ x.\ \neg x !! n by blast+
thus ?thesis unfolding toss-def image-def by auto
qed
```

**lemma** vimage-toss: toss  $n - A = (if \ 1 \in A \ then \ \{s. \ s \ !! \ n\} \ else \ \{\}) \cup (if \ -1 \in A \ then \ \{s. \ \neg s \ !! \ n\} \ else \ \{\})$  **unfolding** vimage-def toss-def **by** auto

**lemma** fortune-Suc: fortune (Suc n) s =fortune n s + toss (Suc n) s **by** (induction n arbitrary: s) (simp add: fortune-def toss-def)+

**lemma** fortune-toss-sum: fortune  $n s = (\sum i \in \{..n\}. toss i s)$ **by** (induction n arbitrary: s) (simp add: fortune-def toss-def, simp add: fortune-Suc)

**lemma** fortune-bound: norm (fortune n s)  $\leq$  Suc n by (induction n) (force simp add: fortune-toss-sum toss-def)+

Our definition of bernoulli-stream constitutes a probability space.

interpretation prob-space bernoulli-stream p unfolding bernoulli-stream-def by (simp add: measure-pmf.prob-space-axioms prob-space.prob-space-stream-space)

**abbreviation** toss-filtration  $p \equiv nat$ -natural-filtration (bernoulli-stream p) toss

The stochastic process toss is adapted to the filtration it generates.

```
(bernoulli-stream p) toss toss
  by (intro nat-adapted-process.intro stochastic-process.adapted-process-natural-filtration)
        (unfold-locales, auto simp add: toss-def bernoulli-stream-def)
Similarly, the stochastic process fortune is adapted to the filtration generated
by the tosses.
interpretation fortune: nat-finite-adapted-process-linorder bernoulli-stream p nat-natural-filtration
(bernoulli-stream p) toss fortune
proof -
    show nat-finite-adapted-process-linorder (bernoulli-stream p) (toss-filtration p)
fortune
   unfolding fortune-toss-sum
   \mathbf{by}\ (intro\ nat\text{-}finite\text{-}adapted\text{-}process\text{-}linorder.intro
                    finite-adapted-process-linorder.intro
                    finite-adapted-process-order.intro
                    finite-adapted-process.intro
                    toss.partial-sum-adapted[folded atMost-atLeast0]) intro-locales
qed
lemma integrable-toss: integrable (bernoulli-stream p) (toss n)
   using toss.random-variable
  by (intro Bochner-Integration.integrable-bound [OF\ integrable-const[of-1::real]])
(auto simp add: toss-def)
lemma integrable-fortune: integrable (bernoulli-stream p) (fortune n) using for-
tune-bound
    by (intro Bochner-Integration.integrable-bound[OF integrable-const[of - Suc n]
fortune.random-variable]) auto
We provide the following lemma to explicitly calculate the probability of
events in this probability space.
{\bf lemma}\ measure-bernoulli-stream-snth-pred:
   assumes 0 \le p and p \le 1 and finite J
   shows prob p \{ w \in space (bernoulli-stream p), \forall j \in J. P j = w !! j \} = p \cap card
(J \cap Collect\ P) * (1 - p) \cap (card\ (J - Collect\ P))
   let ?PiE = (\Pi_E \ i \in J. \ if \ P \ i \ then \ \{True\} \ else \ \{False\})
   have product-prob-space (\lambda-. measure-pmf (bernoulli-pmf p)) by unfold-locales
   hence *: to-stream - '\{s. \ \forall i \in J. \ P \ i = s \ !! \ i\} = \{s. \ \forall i \in J. \ P \ i = s \ i\} using
assms by (simp add: to-stream-def)
   also have ... = prod-emb\ UNIV\ (\lambda-.\ measure-pmf\ (bernoulli-pmf\ p))\ J\ ?PiE
   proof -
          fix s assume (\forall i \in J. P i = s i)
       hence (\forall i \in J. \ P \ i = s \ i) = (s \in prod-emb\ UNIV\ (\lambda -. measure-pmf\ (bernoulli-pmf\ (bernoulli-pmf\
p)) \ J \ ?PiE)
```

interpretation toss: nat-adapted-process bernoulli-stream p nat-natural-filtration

```
by (subst\ prod-emb-iff[of\ s])\ (smt\ (verit,\ best)\ not-def\ assms(3)\ id-def
PiE-eq-singleton UNIV-I extensional-UNIV insert-iff singleton D space-measure-pmf)
        }
        moreover
         {
             fix s assume \neg(\forall i \in J. P i = s i)
             then obtain i where i \in J P i \neq s i by blast
         hence (\forall i \in J. \ P \ i = s \ i) = (s \in prod-emb\ UNIV\ (\lambda -. measure-pmf\ (bernoulli-pmf\ (
p)) \ J \ ?PiE)
               by (simp add: restrict-def prod-emb-iff[of s]) (smt (verit, ccfv-SIG) PiE-mem
assms(3) id-def insert-iff singleton-iff)
        ultimately show ?thesis by auto
    qed
     finally have inteq: (to-stream - '\{s. \forall i \in J. P \ i = s \ !! \ i\}) = prod-emb UNIV
(\lambda-. measure-pmf (bernoulli-pmf p)) J ?PiE.
    let ?M = (Pi_M \ UNIV \ (\lambda -. \ measure-pmf \ (bernoulli-pmf \ p)))
     have emeasure (bernoulli-stream p) \{s \in space \ (bernoulli-stream \ p). \ \forall i \in J. \ P \ i \}
= s !! i  = emeasure ?M (to-stream - '{s. \forall i \in J. P i = s !! i})
          using assms emeasure-distrof to-stream ?M (vimage-algebra (streams (space
(measure-pmf\ (bernoulli-pmf\ p))))\ (!!)\ ?M)\ \{s.\ \forall\ i\in J.\ P\ i=s\ !!\ i\},\ symmetric]
measurable-to-stream[of\ (measure-pmf\ (bernoulli-pmf\ p))]
        by (simp only: bernoulli-stream-def stream-space-def *, simp add: space-PiM)
(smt\ (verit,\ best)\ emeasure-notin-sets\ in-vimage-algebra\ inf-top.right-neutral\ sets-distribution for the sets of the s
vimage-Collect)
    also have ... = emeasure ?M (prod-emb UNIV (\lambda-. measure-pmf (bernoulli-pmf
p)) J ?PiE) using integ by (simp add: space-PiM)
    also have \dots = (\prod i \in J. \text{ emeasure (measure-pmf (bernoulli-pmf p)) (if } P \text{ i then})
\{True\}\ else\ \{False\}))
      by (subst emeasure-PiM-emb) (auto simp add: prob-space-measure-pmf assms(3))
     also have ... = (\prod i \in J \cap Collect \ P. \ ennreal \ p) * (\prod i \in J - Collect \ P. \ ennreal \ p)
      unfolding emeasure-pmf-single of bernoulli-pmf p True, unfolded pmf-bernoulli-True OF
assms(1,2)], symmetric]
                       emeasure-pmf-single[of bernoulli-pmf p False, unfolded pmf-bernoulli-False[OF
assms(1,2)], symmetric]
        by (simp add: prod.Int-Diff[OF assms(3), of - Collect P])
   also have ... = p \ \widehat{} \ card \ (J \cap Collect \ P) * (1 - p) \ \widehat{} \ card \ (J - Collect \ P) using
assms by (simp add: prod-ennreal ennreal-mult' ennreal-power)
      finally show ?thesis using assms by (intro measure-eq-emeasure-eq-ennreal)
auto
qed
lemma
    assumes 0 \le p and p \le 1
    shows measure-bernoulli-stream-snth: prob p \{ w \in space (bernoulli-stream p), w \}
        and measure-bernoulli-stream-neg-snth: prob p \{ w \in space (bernoulli-stream p) \}.
\neg w !! i \} = 1 - p
```

```
Now we can express the expected value of a single coin toss.
lemma integral-toss:
 assumes 0 \le p \ p \le 1
 shows expectation p(toss n) = 2 * p - 1
proof -
 have [simp]:\{s.\ s...\ n\} \in events\ p\ using\ measurable-snth[THEN\ measurable-sets,
of {True} measure-pmf (bernoulli-pmf p) n, folded bernoulli-stream-def]
   by (simp add: vimage-def)
 have expectation \ p \ (toss \ n) = Bochner-Integration.simple-bochner-integral \ (bernoulli-stream
p) (toss n)
   using toss.random-variable[of n, THEN measurable-sets]
  \textbf{by } (intro\ simple-bochner-integrable-eq-integral[symmetric]\ simple-bochner-integrable.intros)
(auto simp add: toss-def simple-function-def image-def)
  also have ... = p - prob p \{s. \neg s !! n\} unfolding simple-bochner-integral-def
using measure-bernoulli-stream-snth[OF assms]
   by (simp add: range-toss, simp add: toss-def)
  also have \dots = p - (1 - prob \ p \ \{s. \ s !! \ n\}) by (subst \ prob-compl[symmetric],
auto simp add: Collect-neg-eq Compl-eq-Diff-UNIV)
 finally show ?thesis using measure-bernoulli-stream-snth[OF assms] by simp
qed
Now, we show that the tosses are independent from one another.
lemma indep-vars-toss:
 assumes 0 \le p \ p \le 1
 shows indep-vars p (\lambda-. borel) toss {\theta..}
proof (subst indep-vars-def, intro conjI indep-sets-sigma)
    fix A J assume asm: J \neq \{\} finite J \forall j \in J. A j \in \{toss \ j - `A \cap space \}
(bernoulli-stream\ p)\ |A.\ A\in borel\}
   hence \forall j \in J. \exists B \in borel. A j = toss j - `B \cap space (bernoulli-stream p) by
   then obtain B where B-is: A j = toss j - B j \cap space (bernoulli-stream p)
B j \in borel \text{ if } j \in J \text{ for } j \text{ by } met is
   have prob p \cap (A \cdot J) = (\prod j \in J. \ prob \ p (A \ j))
   proof cases
We consider the case where there is a zero probability event.
     assume \exists j \in J. 1 \notin B \ j \land -1 \notin B \ j
     then obtain j where j-is: j \in J 1 \notin B j -1 \notin B j by blast
   hence A-j-empty: A j = \{\} using B-is by (force simp add: toss-def vimage-def)
     hence \bigcap (A \cdot J) = \{\} using j-is by blast
     moreover have prob p(A j) = 0 using A-j-empty by simp
     ultimately show ?thesis using j-is asm(2) by auto
   next
```

using measure-bernoulli-stream-snth-pred [OF assms, of  $\{i\}$   $\lambda x$ . True]

measure-bernoulli-stream-snth-pred [OF assms, of  $\{i\}$   $\lambda x$ . False] by auto

```
We now assume all events have positive probability.
          assume \neg(\exists i \in J. \ 1 \notin B \ i \land -1 \notin B \ i)
          hence *: 1 \in B \ j \lor -1 \in B \ j \ \text{if} \ j \in J \ \text{for} \ j \ \text{using} \ that \ \text{by} \ blast
          define J' where [simp]: J' = \{j \in J. (1 \in B \ j) \longleftrightarrow (-1 \notin B \ j)\}
           hence toss \ j \ w \in B \ j \longleftrightarrow (1 \in B \ j) = w \ !! \ j \ if \ j \in J' \ for \ w \ j \ using \ that
unfolding toss-def by simp
          hence (\bigcap (A 'J')) = \{w \in space (bernoulli-stream p). \forall j \in J'. (1 \in B j) = a \in Space (bernoulli-stream p). \forall j \in J'. (1 \in B j) = a \in Space (bernoulli-stream p). \forall j \in J'. (1 \in B j) = a \in Space (bernoulli-stream p). \forall j \in J'. (1 \in B j) = a \in Space (bernoulli-stream p). \forall j \in J'. (1 \in B j) = a \in Space (bernoulli-stream p). \forall j \in J'. (1 \in B j) = a \in Space (bernoulli-stream p). \forall j \in J'. (1 \in B j) = a \in Space (bernoulli-stream p). \forall j \in J'. (1 \in B j) = a \in Space (bernoulli-stream p). \forall j \in J'. (1 \in B j) = a \in Space (bernoulli-stream p). \forall j \in J'. (1 \in B j) = a \in Space (bernoulli-stream p). \forall j \in J'. (1 \in B j) = a \in Space (bernoulli-stream p). \forall j \in J'. (1 \in B j) = a \in Space (bernoulli-stream p). \forall j \in J'. (1 \in B j) = a \in Space (bernoulli-stream p). (1 \in Space (bernoulli-stream p). (2 \in Space (bernoulli-stream p). (3 \in Space (bernoulli-stream p). (4 \in Space (bernoulli-stream p). 
w \parallel j using B-is by force
           hence prob J': prob p ( \cap (A 'J')) = p \cap card (J' \cap \{j. 1 \in B j\}) * (1 - G \cap J')
p) \cap card (J' - \{j, 1 \in B \})
                     using measure-bernoulli-stream-snth-pred[OF assms finite-subset[OF -
asm(2)], of J' \lambda j. 1 \in B j] by auto
The index set J' consists of the indices of all non-trivial events.
          have A-j-True: A j = \{w \in space (bernoulli-stream p), w !! j\} if j \in J' \cap \{j, j\}
1 \in B j for j
              using that by (auto simp add: toss-def B-is(1) split: if-splits)
          have A-j-False: A j = \{w \in space (bernoulli-stream p), \neg w !! j\} if j \in J'
\{j, 1 \in B \}  for j
              using that B-is by (auto simp add: toss-def)
         have A-j-top: A j = space \ (bernoulli-stream \ p) \ \textbf{if} \ j \in J - J' \ \textbf{for} \ j \ \textbf{using} \ that
* by (auto simp add: B-is toss-def)
          hence \bigcap (A ' J) = \bigcap (A ' J') by auto
          hence prob\ p\ (\bigcap\ (A\ `J)) = prob\ p\ (\bigcap\ (A\ `J')) by presburger
           also have ... = (\prod j \in J' \cap \{j, 1 \in B \}), prob p(A j) * (\prod j \in J' - \{j, 1 \in B \})
B j \}. prob p (A j)
            by (simp only: prob-J' A-j-True A-j-False measure-bernoulli-stream-snth[OF]
assms measure-bernoulli-stream-neg-snth[OF assms] conq: prod.conq) simp
       also have ... = (\prod j \in J'. prob p(A j)) using asm(2) by (intro\ prod.Int-Diff[symmetric])
           also have ... = (\prod j \in J'. prob p(A j)) * (\prod j \in J - J'. prob p(A j)) using
A-j-top prob-space by simp
            also have ... = (\prod j \in J. \ prob \ p \ (A \ j)) using asm(2) by (metis \ (no-types,
lifting) J'-def mem-Collect-eq mult.commute prod.subset-diff subsetI)
          finally show ?thesis.
       qed
    thus indep-sets p (\lambda i. {toss i - 'A \cap space (bernoulli-stream p) |A|. A \in sets
borel\) {0..} using measurable-sets[OF toss.random-variable]
   by (intro indep-setsI subsetI) fastforce
qed (simp, intro Int-stableI, simp, metis sets.Int vimage-Int)
```

The fortune of a player is a martingale (resp. sub- or supermartingale) with respect to the filtration generated by the coin tosses.

**theorem** fortune-martingale:

```
assumes p = 1/2
 shows nat-martingale (bernoulli-stream p) (toss-filtration p) fortune
 \mathbf{using}\ cond\text{-}exp\text{-}indep[OF\ for tune. subalg\ indep\text{-}set\text{-}natural\text{-}filtration\ integrable\text{-}toss,}
OF\ zero\text{-}order(1)\ lessI\ indep-vars-toss,\ of\ p
       integral-toss assms
     by (intro fortune.martingale-of-cond-exp-diff-Suc-eq-zero integrable-fortune)
(force simp add: fortune-toss-sum)
theorem fortune-submartingale:
 assumes 1/2 \le p \ p \le 1
 shows nat-submartingale (bernoulli-stream p) (toss-filtration p) fortune
proof (intro fortune.submartingale-of-cond-exp-diff-Suc-nonneg integrable-fortune)
 \mathbf{fix} \ n
 show AE \xi in bernoulli-stream p. 0 \leq cond-exp (bernoulli-stream p) (toss-filtration
(p, n) (\lambda \xi. fortune (Suc n) \xi – fortune (n, \xi) \xi
  using cond-exp-indep[OF fortune.subalq indep-set-natural-filtration integrable-toss,
OF\ zero-order(1)\ lessI\ indep-vars-toss,\ of\ p\ n
         integral-toss[of p Suc n] assms
   by (force simp add: fortune-toss-sum)
qed
theorem fortune-supermartingale:
  assumes 0 \le p \ p \le 1/2
  shows nat-supermartingale (bernoulli-stream p) (toss-filtration p) fortune
proof (intro fortune.supermartingale-of-cond-exp-diff-Suc-le-zero integrable-fortune)
 \mathbf{fix} \ n
 show AE \ \xi in bernoulli-stream p.\ \theta \geq cond-exp (bernoulli-stream p) (toss-filtration
(p, n) (\lambda \xi. fortune (Suc n) \xi – fortune (n, \xi) \xi
  {\bf using} \ cond-exp-indep [OF \ for tune. subalg \ indep-set-natural-filtration \ integrable-toss, \\
OF\ zero\text{-}order(1)\ lessI\ indep-vars-toss,\ of\ p\ n
         integral-toss[of p Suc n] assms
   by (force simp add: fortune-toss-sum)
qed
end
```

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