

On the Formalization of Martingales

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theory	<i>Measure-Space-Addendum</i>	
imports	<i>HOL-Analysis.Measure-Space</i>	
begin		

0.1 Sigma Algebra Generated by a Family of Functions

definition *sigma-gen* :: 'a set \Rightarrow 'b measure \Rightarrow ('a \Rightarrow 'b) set \Rightarrow 'a measure **where**
sigma-gen Ω N $S \equiv \text{sigma } \Omega (\bigcup f \in S. \{f - 'A \cap \Omega \mid A. A \in N\})$

lemma

shows *sets-sigma-gen*: *sets* (*sigma-gen* Ω N S) = *sigma-sets* $\Omega (\bigcup f \in S. \{f - 'A \cap \Omega \mid A. A \in N\})$

and *space-sigma-gen[simp]*: *space* (*sigma-gen* Ω N S) = Ω

by (*auto simp add: sigma-gen-def sets-measure-of-conv space-measure-of-conv*)

lemma *measurable-sigma-gen*:

assumes $f \in S$ $f \in \Omega \rightarrow \text{space } N$

shows $f \in \text{sigma-gen } \Omega$ N $S \rightarrow_M N$

using *assms* **by** (*intro measurableI, auto simp add: sets-sigma-gen*)

lemma *measurable-sigma-gen-singleton*:

assumes $f \in \Omega \rightarrow \text{space } N$

shows $f \in \text{sigma-gen } \Omega$ N $\{f\} \rightarrow_M N$

using *assms* *measurable-sigma-gen* **by** *blast*

lemma *measurable-iff-contains-sigma-gen*:

shows $(f \in M \rightarrow_M N) \longleftrightarrow f \in \text{space } M \rightarrow \text{space } N \wedge \text{sigma-gen } (\text{space } M) N$
 $\{f\} \subseteq M$

proof (*standard, goal-cases*)

case 1

hence $f \in \text{space } M \rightarrow \text{space } N$ **using** *measurable-space* **by** *fast*

thus ?case **unfolding** *sets-sigma-gen* **by** (*simp, intro sigma-algebra.sigma-sets-subset,*
(blast intro: sets.sigma-algebra-axioms measurable-sets[OF 1])+))

next

case 2

thus ?case **using** *measurable-mono*[*OF - refl - space-sigma-gen, of N M*] *measurable-sigma-gen-singleton* **by** *fast*

qed

lemma *measurable-family-iff-contains-sigma-gen*:

shows $(S \subseteq M \rightarrow_M N) \longleftrightarrow S \subseteq \text{space } M \rightarrow \text{space } N \wedge \text{sigma-gen } (\text{space } M)$
 $N S \subseteq M$

proof (*standard, goal-cases*)

case 1

hence *subset*: $S \subseteq \text{space } M \rightarrow \text{space } N$ **using** *measurable-space* **by** *fast*

have $\{f - 'A \cap \text{space } M \mid A. A \in N\} \subseteq M$ **if** $f \in S$ **for** f **using** *measurable-iff-contains-sigma-gen*[*unfolded sets-sigma-gen, of f*] 1 *subset that* **by** *blast*

then show ?case **unfolding** *sets-sigma-gen* **using** *sets.sigma-algebra-axioms* **by**
(simp add: subset, intro sigma-algebra.sigma-sets-subset, blast+))

next

case 2

hence *subset*: $S \subseteq \text{space } M \rightarrow \text{space } N$ **by** *simp*

show ?case

proof (*standard, goal-cases*)

```

    case (1 x)
      have sigma-gen (space M) N {x} ⊆ M by (metis (no-types, lifting) 1 2
sets-sigma-gen SUP-le-iff sigma-sets-le-sets-iff singletonD)
      thus ?case using measurable-iff-contains-sigma-gen subset[THEN subsetD, OF
1] by fast
    qed
  qed

```

```

end
theory Elementary-Metric-Spaces-Addendum
imports HOL-Analysis.Elementary-Metric-Spaces
begin

```

```

lemma diameter-comp-strict-mono:
  fixes s :: nat ⇒ 'a :: metric-space
  assumes strict-mono r bounded {s i | i. r n ≤ i}
  shows diameter {s (r i) | i. n ≤ i} ≤ diameter {s i | i. r n ≤ i}
proof (rule diameter-subset)
  show {s (r i) | i. n ≤ i} ⊆ {s i | i. r n ≤ i} using assms(1) monotoneD
strict-mono-mono by fastforce
qed (intro assms(2))

```

```

lemma diameter-bounded-bound':
  fixes S :: 'a :: metric-space set
  assumes S: bdd-above (case-prod dist ' (S×S)) x ∈ S y ∈ S
  shows dist x y ≤ diameter S
proof -
  have (x,y) ∈ S×S using S by auto
  then have dist x y ≤ (SUP (x,y)∈S×S. dist x y) by (rule cSUP-upper2[OF
assms(1)]) simp
  with ⟨x ∈ S⟩ show ?thesis by (auto simp: diameter-def)
qed

```

```

lemma bounded-imp-dist-bounded:
  assumes bounded (range s)
  shows bounded ((λ(i, j). dist (s i) (s j)) ' ({n..} × {n..}))
  using bounded-dist-comp[OF bounded-fst bounded-snd, OF bounded-Times(1,1)[OF
assms(1,1)]] by (rule bounded-subset, force)

```

```

lemma cauchy-iff-diameter-tends-to-zero-and-bounded:
  fixes s :: nat ⇒ 'a :: metric-space
  shows Cauchy s ⟷ ((λn. diameter {s i | i. i ≥ n}) ⟶ 0 ∧ bounded (range
s))
proof -
  have {s i | i. N ≤ i} ≠ {} for N by blast
  hence diameter-SUP: diameter {s i | i. N ≤ i} = (SUP (i, j) ∈ {N..} × {N..}.
dist (s i) (s j)) for N unfolding diameter-def by (auto intro!: arg-cong[of - - Sup])
  show ?thesis
  proof ((intro iffI) ; clarsimp)

```

```

assume asm: Cauchy s
have  $\exists N. \forall n \geq N. \text{norm } (\text{diameter } \{s\ i \mid i. n \leq i\}) < e$  if e-pos:  $e > 0$  for e
proof –
  obtain N where dist-less:  $\text{dist } (s\ n) (s\ m) < (1/2) * e$  if  $n \geq N\ m \geq N$ 
for n m using asm e-pos by (metis Cauchy-def mult-pos-pos zero-less-divide-iff
zero-less-numeral zero-less-one)
  {
    fix r assume  $r \geq N$ 
    hence  $\text{dist } (s\ n) (s\ m) < (1/2) * e$  if  $n \geq r\ m \geq r$  for n m using dist-less
    that by simp
    hence  $(\text{SUP } (i, j) \in \{r..\} \times \{r..\}. \text{dist } (s\ i) (s\ j)) \leq (1/2) * e$  by (intro
cSup-least) fastforce+
    also have  $\dots < e$  using e-pos by simp
    finally have  $\text{diameter } \{s\ i \mid i. r \leq i\} < e$  using diameter-SUP by presburger
  }
  moreover have  $\text{diameter } \{s\ i \mid i. r \leq i\} \geq 0$  for r unfolding diameter-SUP
using bounded-imp-dist-bounded[OF cauchy-imp-bounded, THEN bounded-imp-bdd-above,
OF asm] by (intro cSup-upper2, auto)
  ultimately show ?thesis by auto
qed
  thus  $(\lambda n. \text{diameter } \{s\ i \mid i. n \leq i\}) \longrightarrow 0 \wedge \text{bounded } (\text{range } s)$  using
cauchy-imp-bounded[OF asm] by (simp add: LIMSEQ-iff)
next
  assume asm:  $(\lambda n. \text{diameter } \{s\ i \mid i. n \leq i\}) \longrightarrow 0$  bounded (range s)
  have  $\exists N. \forall n \geq N. \forall m \geq N. \text{dist } (s\ n) (s\ m) < e$  if e-pos:  $e > 0$  for e
  proof –
    obtain N where diam-less:  $\text{diameter } \{s\ i \mid i. r \leq i\} < e$  if  $r \geq N$  for r
using LIMSEQ-D asm(1) e-pos by fastforce
    {
      fix n m assume  $n \geq N\ m \geq N$ 
      hence  $\text{dist } (s\ n) (s\ m) < e$  using cSUP-lessD[OF bounded-imp-dist-bounded[THEN
bounded-imp-bdd-above], OF asm(2) diam-less[unfolded diameter-SUP]] by auto
    }
    thus ?thesis by blast
  qed
  then show Cauchy s by (simp add: Cauchy-def)
qed
qed

end
theory Bochner-Integration-Addendum
imports HOL–Analysis.Bochner-Integration Elementary-Metric-Spaces-Addendum
begin

```

0.2 Simple Functions

lemma *integrable-implies-simple-function-sequence*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{banach, second-countable-topology}\}$

assumes *integrable M f*
obtains s where $\bigwedge i. \text{simple-function } M (s\ i)$
and $\bigwedge i. \text{emeasure } M \{y \in \text{space } M. s\ i\ y \neq 0\} \neq \infty$
and $\bigwedge x. x \in \text{space } M \implies (\lambda i. s\ i\ x) \longrightarrow f\ x$
and $\bigwedge x\ i. x \in \text{space } M \implies \text{norm } (s\ i\ x) \leq 2 * \text{norm } (f\ x)$
proof –
have $f: f \in \text{borel-measurable } M (\int^+ x. \text{norm } (f\ x) \partial M) < \infty$ **using** *assms*
unfolding *integrable-iff-bounded* **by** *auto*
obtain s where $s: \bigwedge i. \text{simple-function } M (s\ i) \bigwedge x. x \in \text{space } M \implies (\lambda i. s\ i\ x) \longrightarrow f\ x$
 $\bigwedge i\ x. x \in \text{space } M \implies \text{norm } (s\ i\ x) \leq 2 * \text{norm } (f\ x)$ **using**
borel-measurable-implies-sequence-metric[OF f(1)] **unfolding** *norm-conv-dist* **by**
metis
{
fix *i*
have $(\int^+ x. \text{norm } (s\ i\ x) \partial M) \leq (\int^+ x. \text{ennreal } (2 * \text{norm } (f\ x)) \partial M)$ **using**
s **by** (*intro nn-integral-mono, auto*)
also have $\dots < \infty$ **using** *f* **by** (*simp add: nn-integral-cmult ennreal-mult-less-top ennreal-mult*)
finally have *sbi: Bochner-Integration.simple-bochner-integrable M (s i)* **using**
s **by** (*intro simple-bochner-integrableI-bounded*) *auto*
hence $\text{emeasure } M \{y \in \text{space } M. s\ i\ y \neq 0\} \neq \infty$ **by** (*auto intro: integrableI-simple-bochner-integrable simple-bochner-integrable.cases*)
}
thus *?thesis* **using** *that s* **by** *blast*
qed

lemma *simple-function-indicator-representation:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes $f: \text{simple-function } M f$ **and** $x: x \in \text{space } M$
shows $f\ x = (\sum y \in f\ \text{space } M. \text{indicator } (f - \{y\} \cap \text{space } M) x *_R y)$
(is ?l = ?r)

proof –

have $?r = (\sum y \in f\ \text{space } M.$
 $(\text{if } y = f\ x \text{ then indicator } (f - \{y\} \cap \text{space } M) x *_R y \text{ else } 0))$ **by** (*auto intro!: sum.cong*)
also have $\dots = \text{indicator } (f - \{f\ x\} \cap \text{space } M) x *_R f\ x$ **using** *assms* **by** (*auto dest: simple-functionD*)
also have $\dots = f\ x$ **using** *x* **by** (*auto simp: indicator-def*)
finally show *?thesis* **by** *auto*
qed

lemma *simple-function-indicator-representation-AE:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes $f: \text{simple-function } M f$
shows $\text{AE } x \text{ in } M. f\ x = (\sum y \in f\ \text{space } M. \text{indicator } (f - \{y\} \cap \text{space } M) x *_R y)$
by (*metis (mono-tags, lifting) AE-I2 simple-function-indicator-representation f*)

lemmas *simple-function-scaleR[intro] = simple-function-compose2[where h=(*_R)]*

lemmas *integrable-simple-function* = *simple-bochner-integrable.intros*[*THEN has-bochner-integral-simple-bochner.intros*]

lemma *integrable-simple-function-induct*[*consumes 2, case-names cong indicator add, induct set: simple-function*]:

fixes *f* :: 'a \Rightarrow 'b :: {*second-countable-topology, banach*}
assumes *f*: *simple-function* *M f* *emeasure* *M* {*y* \in *space* *M*. *f y* \neq 0} \neq ∞
assumes *cong*: $\bigwedge f g. \text{simple-function } M f \implies \text{emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty$
 $\implies \text{simple-function } M g \implies \text{emeasure } M \{y \in \text{space } M. g y \neq 0\} \neq \infty$

$\implies (\bigwedge x. x \in \text{space } M \implies f x = g x) \implies P f \implies P g$
assumes *indicator*: $\bigwedge A y. A \in \text{sets } M \implies \text{emeasure } M A < \infty \implies P (\lambda x. \text{indicator } A x *_R y)$

assumes *add*: $\bigwedge f g. \text{simple-function } M f \implies \text{emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty \implies$
 $\text{simple-function } M g \implies \text{emeasure } M \{y \in \text{space } M. g y \neq 0\} \neq \infty \implies$

$(\bigwedge z. z \in \text{space } M \implies \text{norm } (f z + g z) = \text{norm } (f z) + \text{norm } (g z)) \implies$
 $P f \implies P g \implies P (\lambda x. f x + g x)$

shows *P f*

proof–

let *?f* = $\lambda x. (\sum y \in f \text{ 'space } M. \text{indicat-real } (f - \text{' } \{y\} \cap \text{space } M) x *_R y)$
have *f-ae-eq*: *f x* = *?f x* **if** *x* \in *space* *M* **for** *x* **using** *simple-function-indicator-representation*[*OF f(1) that*].

moreover have *emeasure* *M* {*y* \in *space* *M*. *?f y* \neq 0} \neq ∞ **by** (*metis* (*no-types, lifting*) *Collect-cong* *calculation f(2)*)

moreover have *P* ($\lambda x. \sum y \in S. \text{indicat-real } (f - \text{' } \{y\} \cap \text{space } M) x *_R y$)
 $\text{simple-function } M (\lambda x. \sum y \in S. \text{indicat-real } (f - \text{' } \{y\} \cap \text{space } M) x *_R y)$

emeasure *M* {*y* \in *space* *M*. $(\sum x \in S. \text{indicat-real } (f - \text{' } \{x\} \cap \text{space } M) y *_R x) \neq 0$ } \neq ∞

if *S* \subseteq *f* ' *space* *M* **for** *S* **using** *simple-functionD(1)*[*OF assms(1), THEN rev-finite-subset, OF that*] **that**

proof (*induction rule: finite-induct*)

case *empty*

{

case 1

then show *?case* **using** *indicator*[*of {}*] **by force**

next

case 2

then show *?case* **by force**

next

case 3

then show *?case* **by force**

}

next
case (*insert* x F)
have $(f - \{x\} \cap \text{space } M) \subseteq \{y \in \text{space } M. f y \neq 0\}$ **if** $x \neq 0$ **using** *that* **by** *blast*
moreover **have** $\{y \in \text{space } M. f y \neq 0\} = \text{space } M - (f - \{0\} \cap \text{space } M)$
by *blast*
moreover **have** $\text{space } M - (f - \{0\} \cap \text{space } M) \in \text{sets } M$ **using** *simple-functionD(2)[OF f(1)]* **by** *blast*
ultimately **have** *fin-0*: $\text{emeasure } M (f - \{x\} \cap \text{space } M) < \infty$ **if** $x \neq 0$
using *that* **by** (*metis emeasure-mono f(2) infinity-enreal-def top.not-eq-extremum top-unique*)
hence *fin-1*: $\text{emeasure } M \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\} \neq \infty$ **if** $x \neq 0$ **by** (*metis (mono-tags, lifting) emeasure-mono f(1) indicator-simps(2) linorder-not-less mem-Collect-eq scaleR-eq-0-iff simple-functionD(2) subsetI* *that*)

have *: $(\sum y \in \text{insert } x F. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y) = (\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y) + \text{indicat-real } (f - \{x\} \cap \text{space } M) x *_R x$ **for** x **by** (*metis (no-types, lifting) Diff-empty Diff-insert0 add.commute insert.hyps(1) insert.hyps(2) sum.insert-remove*)
have **: $\{y \in \text{space } M. (\sum x \in \text{insert } x F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \subseteq \{y \in \text{space } M. (\sum x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\}$ **unfolding** * **by** *fastforce*
{
case 1
hence $x: x \in f - \text{space } M$ **and** $F: F \subseteq f - \text{space } M$ **by** *auto*
show ?*case*
proof (*cases* $x = 0$)
case *True*
then **show** ?*thesis* **unfolding** * **using** *insert(3)[OF F]* **by** *simp*
next
case *False*
have *norm-argument*: $\text{norm } ((\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) + \text{indicat-real } (f - \{x\} \cap \text{space } M) z *_R x) = \text{norm } (\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) + \text{norm } (\text{indicat-real } (f - \{x\} \cap \text{space } M) z *_R x)$
if $z: z \in \text{space } M$ **for** z
proof (*cases* $f z = x$)
case *True*
have $\text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y = 0$ **if** $y \in F$ **for** y **using** *True insert(2) z that 1 unfolding indicator-def* **by** *force*
hence $(\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) = 0$ **by** (*meson sum.neutral*)
then **show** ?*thesis* **by** *force*
next
case *False*
then **show** ?*thesis* **by** *force*
qed
show ?*thesis* **using** *False simple-functionD(2)[OF f(1)] insert(3,5)[OF F]*

```

simple-function-scaleR fin-0 fin-1 by (subst *, subst add, subst simple-function-sum)
(blast intro: norm-argument indicator)+
  qed
next
  case 2
  hence x:  $x \in f \text{ ' space } M$  and  $F: F \subseteq f \text{ ' space } M$  by auto
  show ?case
  proof (cases  $x = 0$ )
    case True
    then show ?thesis unfolding * using insert(4)[OF F] by simp
  next
    case False
    then show ?thesis unfolding * using insert(4)[OF F] simple-functionD(2)[OF
f(1)] by fast
  qed
next
  case 3
  hence x:  $x \in f \text{ ' space } M$  and  $F: F \subseteq f \text{ ' space } M$  by auto
  show ?case
  proof (cases  $x = 0$ )
    case True
    then show ?thesis unfolding * using insert(5)[OF F] by simp
  next
    case False
    have emeasure M  $\{y \in \text{space } M. (\sum_{x \in \text{insert } x F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \leq \text{emeasure } M \{y \in \text{space } M. (\sum_{x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\}$ 
    using ** simple-functionD(2)[OF insert(4)[OF F]] simple-functionD(2)[OF
f(1)] by (intro emeasure-mono, force+)
    also have ...  $\leq \text{emeasure } M \{y \in \text{space } M. (\sum_{x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} + \text{emeasure } M \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\}$ 
    using simple-functionD(2)[OF insert(4)[OF F]] simple-functionD(2)[OF
f(1)] by (intro emeasure-subadditive, force+)
    also have ...  $< \infty$  using insert(5)[OF F] fin-1[OF False] by (simp add:
less-top)
    finally show ?thesis by simp
  qed
}
qed
moreover have simple-function M  $(\lambda x. \sum_{y \in f \text{ ' space } M. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y)$  using calculation by blast
moreover have P  $(\lambda x. \sum_{y \in f \text{ ' space } M. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y)$  using calculation by blast
ultimately show ?thesis by (intro cong[OF - - f(1,2)], blast, presburger+)
qed

```


lemma *integrable-simple-function-induct-nn*[consumes 3, case-names *cong indicator add, induct set: simple-function*]:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes f : *simple-function* M f *emeasure* M $\{y \in \text{space } M. f\ y \neq 0\} \neq \infty \wedge x. x \in \text{space } M \implies f\ x \geq 0$

assumes *cong*: $\wedge f\ g. \text{simple-function } M\ f \implies \text{emeasure } M\ \{y \in \text{space } M. f\ y \neq 0\} \neq \infty \implies (\wedge x. x \in \text{space } M \implies f\ x \geq 0) \implies \text{simple-function } M\ g \implies \text{emeasure } M\ \{y \in \text{space } M. g\ y \neq 0\} \neq \infty \implies (\wedge x. x \in \text{space } M \implies g\ x \geq 0) \implies (\wedge x. x \in \text{space } M \implies f\ x = g\ x) \implies P\ f \implies P\ g$

assumes *indicator*: $\wedge A\ y. y \geq 0 \implies A \in \text{sets } M \implies \text{emeasure } M\ A < \infty \implies P\ (\lambda x. \text{indicator } A\ x\ *_{\mathbb{R}}\ y)$

assumes *add*: $\wedge f\ g. (\wedge x. x \in \text{space } M \implies f\ x \geq 0) \implies \text{simple-function } M\ f \implies \text{emeasure } M\ \{y \in \text{space } M. f\ y \neq 0\} \neq \infty \implies$

$(\wedge x. x \in \text{space } M \implies g\ x \geq 0) \implies \text{simple-function } M\ g \implies \text{emeasure } M\ \{y \in \text{space } M. g\ y \neq 0\} \neq \infty \implies$

$(\wedge z. z \in \text{space } M \implies \text{norm } (f\ z + g\ z) = \text{norm } (f\ z) + \text{norm } (g\ z)) \implies$

$P\ f \implies P\ g \implies P\ (\lambda x. f\ x + g\ x)$

shows $P\ f$

proof–

let $?f = \lambda x. (\sum y \in f\ \text{'space } M. \text{indicat-real } (f - \text{'}\{y\} \cap \text{space } M)\ x\ *_{\mathbb{R}}\ y)$

have $f\text{-ae-eq}$: $f\ x = ?f\ x$ **if** $x \in \text{space } M$ **for** x **using** *simple-function-indicator-representation*[*OF* $f(1)$ *that*].

moreover have *emeasure* $M\ \{y \in \text{space } M. ?f\ y \neq 0\} \neq \infty$ **by** (*metis* (*no-types, lifting*) *Collect-cong calculation f(2)*)

moreover have $P\ (\lambda x. \sum y \in S. \text{indicat-real } (f - \text{'}\{y\} \cap \text{space } M)\ x\ *_{\mathbb{R}}\ y)$
 $\text{simple-function } M\ (\lambda x. \sum y \in S. \text{indicat-real } (f - \text{'}\{y\} \cap \text{space } M)\ x\ *_{\mathbb{R}}\ y)$

emeasure $M\ \{y \in \text{space } M. (\sum x \in S. \text{indicat-real } (f - \text{'}\{x\} \cap \text{space } M)\ y\ *_{\mathbb{R}}\ x) \neq 0\} \neq \infty$

$\wedge x. x \in \text{space } M \implies 0 \leq (\sum y \in S. \text{indicat-real } (f - \text{'}\{y\} \cap \text{space } M)\ x\ *_{\mathbb{R}}\ y)$

if $S \subseteq f\ \text{'space } M$ **for** S **using** *simple-functionD(1)*[*OF* *assms(1)*, *THEN rev-finite-subset, OF that*] *that*

proof (*induction rule: finite-subset-induct'*)

case *empty*

{

case 1

then show $?case$ **using** *indicator*[*of* 0 { $\}$] **by force**

next

case 2

then show $?case$ **by force**

next

case 3

then show $?case$ **by force**

next

case 4

then show $?case$ **by force**

```

}
next
  case (insert x F)
  have (f - ' {x} ∩ space M) ⊆ {y ∈ space M. f y ≠ 0} if x ≠ 0 using that by
blast
  moreover have {y ∈ space M. f y ≠ 0} = space M - (f - ' {0} ∩ space M)
by blast
  moreover have space M - (f - ' {0} ∩ space M) ∈ sets M using simple-functionD(2)[OF f(1)] by blast
  ultimately have fin-0: emeasure M (f - ' {x} ∩ space M) < ∞ if x ≠ 0
using that by (metis emeasure-mono f(2) infinity-ennreal-def top.not-eq-extremum
top-unique)
  hence fin-1: emeasure M {y ∈ space M. indicat-real (f - ' {x} ∩ space M) y *R
x ≠ 0} ≠ ∞ if x ≠ 0 by (metis (mono-tags, lifting) emeasure-mono f(1) indica-
tor-simps(2) linorder-not-less mem-Collect-eq scaleR-eq-0-iff simple-functionD(2)
subsetI that)

  have nonneg-x: x ≥ 0 using insert f by blast
  have *: (∑ y ∈ insert x F. indicat-real (f - ' {y} ∩ space M) xa *R y) =
(∑ y ∈ F. indicat-real (f - ' {y} ∩ space M) xa *R y) + indicat-real (f - ' {x} ∩
space M) xa *R x for xa by (metis (no-types, lifting) add commute insert.hyps(1)
insert.hyps(4) sum.insert)
  have **: {y ∈ space M. (∑ x ∈ insert x F. indicat-real (f - ' {x} ∩ space M) y
*R x) ≠ 0} ⊆ {y ∈ space M. (∑ x ∈ F. indicat-real (f - ' {x} ∩ space M) y *R x)
≠ 0} ∪ {y ∈ space M. indicat-real (f - ' {x} ∩ space M) y *R x ≠ 0} unfolding
* by fastforce
  {
    case 1
    show ?case
    proof (cases x = 0)
      case True
      then show ?thesis unfolding * using insert by simp
    next
      case False
      have norm-argument: norm ((∑ y ∈ F. indicat-real (f - ' {y} ∩ space M) z
*R y) + indicat-real (f - ' {x} ∩ space M) z *R x) = norm (∑ y ∈ F. indicat-real
(f - ' {y} ∩ space M) z *R y) + norm (indicat-real (f - ' {x} ∩ space M) z *R x)
if z: z ∈ space M for z
      proof (cases f z = x)
        case True
        have indicat-real (f - ' {y} ∩ space M) z *R y = 0 if y ∈ F for y using
True insert z that 1 unfolding indicator-def by force
        hence (∑ y ∈ F. indicat-real (f - ' {y} ∩ space M) z *R y) = 0 by (meson
sum.neutral)
        thus ?thesis by force
      qed (force)
      show ?thesis using False fin-0 fin-1 f norm-argument by (subst *, subst add,
presburger add: insert, intro insert, intro insert, fastforce simp add: indicator-def
intro!: insert(2) f(3), auto intro!: indicator insert nonneg-x)

```

```

    qed
  next
    case 2
    show ?case
    proof (cases x = 0)
      case True
      then show ?thesis unfolding * using insert by simp
    next
      case False
      then show ?thesis unfolding * using insert simple-functionD(2)[OF f(1)]
    by fast
    qed
  next
    case 3
    show ?case
    proof (cases x = 0)
      case True
      then show ?thesis unfolding * using insert by simp
    next
      case False
      have emeasure M {y ∈ space M. (∑ x ∈ insert x F. indicat-real (f - ' {x}
        ∩ space M) y *R x) ≠ 0} ≤ emeasure M ({y ∈ space M. (∑ x ∈ F. indicat-real (f
        - ' {x} ∩ space M) y *R x) ≠ 0} ∪ {y ∈ space M. indicat-real (f - ' {x} ∩ space
        M) y *R x ≠ 0})
      using ** simple-functionD(2)[OF insert(6)] simple-functionD(2)[OF f(1)]
      insert.IH(2) by (intro emeasure-mono, blast, simp)
      also have ... ≤ emeasure M {y ∈ space M. (∑ x ∈ F. indicat-real (f - ' {x}
        ∩ space M) y *R x) ≠ 0} + emeasure M {y ∈ space M. indicat-real (f - ' {x} ∩
        space M) y *R x ≠ 0}
      using simple-functionD(2)[OF insert(6)] simple-functionD(2)[OF f(1)]
    by (intro emeasure-subadditive, force+)
      also have ... < ∞ using insert(7) fin-1[OF False] by (simp add: less-top)
      finally show ?thesis by simp
    qed
  next
    case (4 ξ)
    thus ?case using insert nonneg-x f(3) by (auto simp add: scaleR-nonneg-nonneg
    intro: sum-nonneg)
  }
  qed
  moreover have simple-function M (λx. ∑ y ∈ f - ' space M. indicat-real (f - ' {y}
    ∩ space M) x *R y) using calculation by blast
  moreover have P (λx. ∑ y ∈ f - ' space M. indicat-real (f - ' {y} ∩ space M) x
    *R y) using calculation by blast
  moreover have ∧ x. x ∈ space M ⇒ 0 ≤ f x using f(3) by simp
  ultimately show ?thesis by (intro cong[OF - - - f(1,2)], blast, blast, fast)
presburger+
qed

```

lemma *finite-nn-integral-imp-ae-finite*:
fixes $f :: 'a \Rightarrow \text{ennreal}$
assumes $f \in \text{borel-measurable } M \ (\int^{+x}. f \ x \ \partial M) < \infty$
shows $AE \ x \text{ in } M. f \ x < \infty$
proof (rule *ccontr*, goal-cases)
case 1
let $?A = \text{space } M \cap \{x. f \ x = \infty\}$
have $*$: $\text{emeasure } M \ ?A > 0$ **using** 1 **assms**(1) **by** (metis (mono-tags, lifting)
assms(2) *eventually-mono infinity-ennreal-def nn-integral-noteq-infinite top.not-eq-extremum*)
have $(\int^{+x}. f \ x * \text{indicator } ?A \ x \ \partial M) = (\int^{+x}. \infty * \text{indicator } ?A \ x \ \partial M)$ **by**
(metis (mono-tags, lifting) *indicator-inter-arith indicator-simps*(2) *mem-Collect-eq*
mult-eq-0-iff)
also have $\dots = \infty * \text{emeasure } M \ ?A$ **using** *assms*(1) **by** (intro *nn-integral-cmult-indicator*,
simp)
also have $\dots = \infty$ **using** $*$ **by** *fastforce*
finally have $(\int^{+x}. f \ x * \text{indicator } ?A \ x \ \partial M) = \infty$.
moreover have $(\int^{+x}. f \ x * \text{indicator } ?A \ x \ \partial M) \leq (\int^{+x}. f \ x \ \partial M)$ **by** (intro
nn-integral-mono, *simp add: indicator-def*)
ultimately show $?case$ **using** *assms*(2) **by** *simp*
qed

lemma *cauchy-L1-AE-cauchy-subseq*:
fixes $s :: \text{nat} \Rightarrow 'a \Rightarrow 'b :: \{\text{banach, second-countable-topology}\}$
assumes [*measurable*]: $\bigwedge n. \text{integrable } M \ (s \ n)$
and $\bigwedge e. e > 0 \implies \exists N. \forall i \geq N. \forall j \geq N. \text{LINT } x | M. \text{norm } (s \ i \ x - s \ j \ x) < e$
obtains r **where** *strict-mono* $r \ AE \ x \text{ in } M. \text{Cauchy } (\lambda i. s \ (r \ i) \ x)$
proof–
have $\exists r. \forall n. (\forall i \geq r \ n. \forall j \geq r \ n. \text{LINT } x | M. \text{norm } (s \ i \ x - s \ j \ x) < (1 / 2) ^ n) \wedge (r \ (\text{Suc } n) > r \ n)$
proof (intro *dependent-nat-choice*, goal-cases)
case 1
then show $?case$ **using** *assms*(2) **by** *presburger*
next
case (2 $x \ n$)
obtain N **where** $*$: $\text{LINT } x | M. \text{norm } (s \ i \ x - s \ j \ x) < (1 / 2) ^ \text{Suc } n$ **if** $i \geq N$
for $i \ j$ **using** *assms*(2)[*of* $(1 / 2) ^ \text{Suc } n$] **by** *auto*
{
fix $i \ j$ **assume** $i \geq \max N \ (\text{Suc } x) \ j \geq \max N \ (\text{Suc } x)$
hence $\text{integral}^L M \ (\lambda x. \text{norm } (s \ i \ x - s \ j \ x)) < (1 / 2) ^ \text{Suc } n$ **using** $*$ **by**
fastforce
}
then show $?case$ **by** *fastforce*
qed
then obtain r **where** *strict-mono*: *strict-mono* r **and** $\forall i \geq r \ n. \forall j \geq r \ n. \text{LINT } x | M. \text{norm } (s \ i \ x - s \ j \ x) < (1 / 2) ^ n$ **for** n **using** *strict-mono-Suc-iff* **by** *blast*
hence r -is: $\text{LINT } x | M. \text{norm } (s \ (r \ (\text{Suc } n)) \ x - s \ (r \ n) \ x) < (1 / 2) ^ n$ **for** n

by (*simp add: strict-mono-leD*)

define *g* **where** $g = (\lambda n\ x.\ (\sum i \leq n.\ \text{ennreal}\ (\text{norm}\ (s\ (r\ (\text{Suc}\ i))\ x - s\ (r\ i)\ x))))$

define *g'* **where** $g' = (\lambda x.\ \sum i.\ \text{ennreal}\ (\text{norm}\ (s\ (r\ (\text{Suc}\ i))\ x - s\ (r\ i)\ x)))$

have *integrable-g*: $(\int^+ x.\ g\ n\ x\ \partial M) < 2$ **for** *n*

proof –

have $(\int^+ x.\ g\ n\ x\ \partial M) = (\int^+ x.\ (\sum i \leq n.\ \text{ennreal}\ (\text{norm}\ (s\ (r\ (\text{Suc}\ i))\ x - s\ (r\ i)\ x)))\ \partial M)$ **using** *g-def* **by** *simp*

also have $\dots = (\sum i \leq n.\ (\int^+ x.\ \text{ennreal}\ (\text{norm}\ (s\ (r\ (\text{Suc}\ i))\ x - s\ (r\ i)\ x))\ \partial M))$ **by** (*intro nn-integral-sum, simp*)

also have $\dots = (\sum i \leq n.\ \text{LINT } x | M.\ \text{norm}\ (s\ (r\ (\text{Suc}\ i))\ x - s\ (r\ i)\ x))$ **unfolding** *dist-norm* **using** *assms(1)* **by** (*subst nn-integral-eq-integral[OF integrable-norm], auto*)

also have $\dots < \text{ennreal}\ (\sum i \leq n.\ (1 / 2) ^ i)$ **by** (*intro ennreal-lessI[OF sum-pos sum-strict-mono[OF finite-atMost - r-is]], auto*)

also have $\dots \leq \text{ennreal } 2$ **unfolding** *sum-gp0[of 1 / 2 n]* **by** (*intro ennreal-leI, auto*)

finally show $(\int^+ x.\ g\ n\ x\ \partial M) < 2$ **by** *simp*

qed

have *integrable-g'*: $(\int^+ x.\ g'\ x\ \partial M) \leq 2$

proof –

have *incseq* $(\lambda n.\ g\ n\ x)$ **for** *x* **by** (*intro incseq-SucI, auto simp add: g-def ennreal-leI*)

hence *convergent* $(\lambda n.\ g\ n\ x)$ **for** *x* **unfolding** *convergent-def* **using** *LIM-SEQ-incseq-SUP* **by** *fast*

hence $(\lambda n.\ g\ n\ x) \longrightarrow g'\ x$ **for** *x* **unfolding** *g-def g'-def* **by** (*intro summable-iff-convergent'[THEN iffD2, THEN summable-LIMSEQ], blast*)

hence $(\int^+ x.\ g'\ x\ \partial M) = (\int^+ x.\ \liminf (\lambda n.\ g\ n\ x)\ \partial M)$ **by** (*metis lim-imp-Liminf trivial-limit-sequentially*)

also have $\dots \leq \liminf (\lambda n.\ \int^+ x.\ g\ n\ x\ \partial M)$ **by** (*intro nn-integral-liminf, simp add: g-def*)

also have $\dots \leq \liminf (\lambda n.\ 2)$ **using** *integrable-g* **by** (*intro Liminf-mono*) (*simp add: order-le-less*)

also have $\dots = 2$ **using** *sequentially-bot tendsto-iff-Liminf-eq-Limsup* **by** *blast*

finally show *?thesis* .

qed

hence *AE* *x* *in* *M*. $g'\ x < \infty$ **by** (*intro finite-nn-integral-imp-ae-finite*) (*auto simp add: order-le-less-trans g'-def*)

moreover have *summable* $(\lambda i.\ \text{norm}\ (s\ (r\ (\text{Suc}\ i))\ x - s\ (r\ i)\ x))$ **if** $g'\ x \neq \infty$ **for** *x* **using** *that* **unfolding** *g'-def* **by** (*intro summable-suminf-not-top*) *fastforce*+

ultimately have *ae-summable*: *AE* *x* *in* *M*. *summable* $(\lambda i.\ s\ (r\ (\text{Suc}\ i))\ x - s\ (r\ i)\ x)$ **using** *summable-norm-cancel* **unfolding** *dist-norm* **by** *force*

{
 fix *x* **assume** *summable* $(\lambda i.\ s\ (r\ (\text{Suc}\ i))\ x - s\ (r\ i)\ x)$

hence $(\lambda n. \sum i < n. s(r(Suc\ i))\ x - s(r\ i)\ x) \longrightarrow (\sum i. s(r(Suc\ i))\ x - s(r\ i)\ x)$ **using** *summable-LIMSEQ* **by** *blast*
moreover have $(\lambda n. (\sum i < n. s(r(Suc\ i))\ x - s(r\ i)\ x)) = (\lambda n. s(r\ n)\ x - s(r\ 0)\ x)$ **using** *sum-lessThan-telescope* **by** *fastforce*
ultimately have $(\lambda n. s(r\ n)\ x - s(r\ 0)\ x) \longrightarrow (\sum i. s(r(Suc\ i))\ x - s(r\ i)\ x)$ **by** *argo*
hence $(\lambda n. s(r\ n)\ x - s(r\ 0)\ x + s(r\ 0)\ x) \longrightarrow (\sum i. s(r(Suc\ i))\ x - s(r\ i)\ x) + s(r\ 0)\ x$ **by** *(intro isCont-tendsto-compose[of - $\lambda z. z + s(r\ 0)\ x$], auto)*
hence *Cauchy* $(\lambda n. s(r\ n)\ x)$ **by** *(simp add: LIMSEQ-imp-Cauchy)*
}
hence *AE* x *in* M . *Cauchy* $(\lambda i. s(r\ i)\ x)$ **using** *ae-summable* **by** *fast*
thus *?thesis* **by** *(rule that[OF strict-mono(1)])*
qed

0.3 Linearly Ordered Banach Spaces

lemma *integrable-max[simp, intro]:*
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology}\}$
assumes $fg[\text{measurable}]: \text{integrable } M\ f\ \text{integrable } M\ g$
shows *integrable* $M\ (\lambda x. \max(f\ x)\ (g\ x))$
proof *(rule Bochner-Integration.integrable-bound)*
{
fix $x\ y :: 'b$
have $\text{norm } (\max\ x\ y) \leq \max(\text{norm } x)\ (\text{norm } y)$ **by** *linarith*
also have $\dots \leq \text{norm } x + \text{norm } y$ **by** *simp*
finally have $\text{norm } (\max\ x\ y) \leq \text{norm } (\text{norm } x + \text{norm } y)$ **by** *simp*
}
thus *AE* x *in* M . $\text{norm } (\max(f\ x)\ (g\ x)) \leq \text{norm } (\text{norm } (f\ x) + \text{norm } (g\ x))$ **by** *simp*
qed *(auto intro: Bochner-Integration.integrable-add[OF integrable-norm[OF fg(1)] integrable-norm[OF fg(2)])*

lemma *integrable-min[simp, intro]:*
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology}\}$
assumes $[measurable]: \text{integrable } M\ f\ \text{integrable } M\ g$
shows *integrable* $M\ (\lambda x. \min(f\ x)\ (g\ x))$
proof *-*
have $\text{norm } (\min(f\ x)\ (g\ x)) \leq \text{norm } (f\ x) \vee \text{norm } (\min(f\ x)\ (g\ x)) \leq \text{norm } (g\ x)$ **for** x **by** *linarith*
thus *?thesis* **by** *(intro integrable-bound[OF integrable-max[OF integrable-norm(1,1), OF assms]], auto)*
qed

lemma *integral-nonneg-AE-banach:*
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes $[measurable]: f \in \text{borel-measurable } M$ **and** *nonneg: AE* x *in* M . $0 \leq f\ x$
shows $0 \leq \text{integral}^L\ M\ f$
proof *cases*

assume *integrable*: *integrable* $M f$
hence \max : $(\lambda x. \max 0 (f x)) \in \text{borel-measurable } M \wedge x. 0 \leq \max 0 (f x)$
integrable $M (\lambda x. \max 0 (f x))$ **by** *auto*
hence $0 \leq \text{integral}^L M (\lambda x. \max 0 (f x))$
proof –
obtain s **where** $*$: $\wedge i. \text{simple-function } M (s i)$
 $\wedge i. \text{emeasure } M \{y \in \text{space } M. s i y \neq 0\} \neq \infty$
 $\wedge x. x \in \text{space } M \implies (\lambda i. s i x) \longrightarrow f x$
 $\wedge x i. x \in \text{space } M \implies \text{norm } (s i x) \leq 2 * \text{norm } (f x)$ **using**
integrable-implies-simple-function-sequence[*OF integrable*] **by** *blast*
have *simple*: $\wedge i. \text{simple-function } M (\lambda x. \max 0 (s i x))$ **using** $*$ **by** *fast*
have $\wedge i. \{y \in \text{space } M. \max 0 (s i y) \neq 0\} \subseteq \{y \in \text{space } M. s i y \neq 0\}$
unfolding *max-def* **by** *force*
moreover **have** $\wedge i. \{y \in \text{space } M. s i y \neq 0\} \in \text{sets } M$ **using** $*$ **by** *measurable*
ultimately **have** $\wedge i. \text{emeasure } M \{y \in \text{space } M. \max 0 (s i y) \neq 0\} \leq$
 $\text{emeasure } M \{y \in \text{space } M. s i y \neq 0\}$ **by** (*rule emeasure-mono*)
hence $**$: $\wedge i. \text{emeasure } M \{y \in \text{space } M. \max 0 (s i y) \neq 0\} \neq \infty$ **using** $*(2)$
by (*auto intro: order.strict-trans1 simp add: top.not-eq-extremum*)
have $\wedge x. x \in \text{space } M \implies (\lambda i. \max 0 (s i x)) \longrightarrow \max 0 (f x)$ **using** $*(3)$
tendsto-max **by** *blast*
moreover **have** $\wedge x i. x \in \text{space } M \implies \text{norm } (\max 0 (s i x)) \leq \text{norm } (2 *_{\mathbb{R}}$
 $f x)$ **using** $*(4)$ **unfolding** *max-def* **by** *auto*
ultimately **have** *tendsto*: $(\lambda i. \text{integral}^L M (\lambda x. \max 0 (s i x))) \longrightarrow \text{integral}^L$
 $M (\lambda x. \max 0 (f x))$
using *borel-measurable-simple-function simple integrable* **by** (*intro inte-*
gral-dominated-convergence[*OF max(1) - integrable-norm*[*OF integrable-scaleR-right*],
of - 2 f], *blast+*)
{
fix $h :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$
dered-real-vector\}
assume *simple-function* $M h$ *emeasure* $M \{y \in \text{space } M. h y \neq 0\} \neq \infty \wedge x.$
 $x \in \text{space } M \longrightarrow h x \geq 0$
hence $*$: $\text{integral}^L M h \geq 0$
proof (*induct rule: integrable-simple-function-induct-nn*)
case (*cong f g*)
then show *?case* **using** *Bochner-Integration.integral-cong* **by** *force*
next
case (*indicator A y*)
hence $A \neq \{\}$ $\implies y \geq 0$ **using** *sets.sets-into-space* **by** *fastforce*
then show *?case* **using** *indicator* **by** (*cases A = \{\}*, *auto simp add:*
scaleR-nonneg-nonneg)
next
case (*add f g*)
then show *?case* **by** (*simp add: integrable-simple-function*)
qed
}
thus *?thesis* **using** *LIMSEQ-le-const*[*OF tendsto, of 0*] $** \text{simple}$ **by** *fastforce*
qed
also **have** $\dots = \text{integral}^L M f$ **using** *nonneg* **by** (*auto intro: integral-cong-AE*)

finally show *?thesis* .
qed (*simp add: not-integrable-integral-eq*)

lemma *integral-mono-AE-banach*:
 fixes $f\ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
 assumes *integrable* $M\ f$ *integrable* $M\ g$ *AE* x in M . $f\ x \leq g\ x$
 shows $\text{integral}^L\ M\ f \leq \text{integral}^L\ M\ g$
 using *integral-nonneg-AE-banach*[*of* $\lambda x. g\ x - f\ x$] *assms* *Bochner-Integration.integral-diff*[*OF* *assms*(1,2)] **by** *force*

lemma *integral-mono-banach*:
 fixes $f\ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
 assumes *integrable* $M\ f$ *integrable* $M\ g$ $\bigwedge x. x \in \text{space } M \implies f\ x \leq g\ x$
 shows $\text{integral}^L\ M\ f \leq \text{integral}^L\ M\ g$
 using *integral-mono-AE-banach* *assms* **by** *blast*

0.4 Integrability and Measurability of the Diameter

context
 fixes $s :: \text{nat} \Rightarrow 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$ **and** M
 assumes *bounded*: $\bigwedge x. x \in \text{space } M \implies \text{bounded } (\text{range } (\lambda i. s\ i\ x))$
begin

lemma *borel-measurable-diameter*:
 assumes [*measurable*]: $\bigwedge i. (s\ i) \in \text{borel-measurable } M$
 shows $(\lambda x. \text{diameter } \{s\ i\ x \mid i. n \leq i\}) \in \text{borel-measurable } M$
proof –
 have $\{s\ i\ x \mid i. N \leq i\} \neq \{\}$ **for** $x\ N$ **by** *blast*
 hence *diameter-SUP*: $\text{diameter } \{s\ i\ x \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N.. \} \times \{N.. \}. \text{dist } (s\ i\ x) (s\ j\ x))$ **for** $x\ N$ **unfolding** *diameter-def* **by** (*auto intro!: arg-cong[of - - Sup]*)

 have *case-prod dist* ‘ $(\{s\ i\ x \mid i. n \leq i\} \times \{s\ i\ x \mid i. n \leq i\}) = ((\lambda(i, j). \text{dist } (s\ i\ x) (s\ j\ x))) '(\{n.. \} \times \{n.. \}))$ **for** x **by** *fast*
 hence *: $(\lambda x. \text{diameter } \{s\ i\ x \mid i. n \leq i\}) = (\lambda x. \text{Sup } ((\lambda(i, j). \text{dist } (s\ i\ x) (s\ j\ x))) '(\{n.. \} \times \{n.. \})))$ **using** *diameter-SUP* **by** (*simp add: case-prod-beta'*)

 have *bounded* $((\lambda(i, j). \text{dist } (s\ i\ x) (s\ j\ x))) '(\{n.. \} \times \{n.. \}))$ **if** $x \in \text{space } M$ **for** x **by** (*rule bounded-imp-dist-bounded*[*OF* *bounded, OF that*])
 hence *bdd*: *bdd-above* $((\lambda(i, j). \text{dist } (s\ i\ x) (s\ j\ x))) '(\{n.. \} \times \{n.. \}))$ **if** $x \in \text{space } M$ **for** x **using** *that bounded-imp-bdd-above* **by** *presburger*
 have *fst* $p \in \text{borel-measurable } M$ *snd* $p \in \text{borel-measurable } M$ **if** $p \in s ' \{n.. \} \times s ' \{n.. \}$ **for** p **using** *that* **by** *fastforce+*
 hence $(\lambda x. \text{fst } p\ x - \text{snd } p\ x) \in \text{borel-measurable } M$ **if** $p \in s ' \{n.. \} \times s ' \{n.. \}$ **for** p **using** *that borel-measurable-diff* **by** *simp*
 hence $(\lambda x. \text{case } p \text{ of } (f, g) \Rightarrow \text{dist } (f\ x) (g\ x)) \in \text{borel-measurable } M$ **if** $p \in s ' \{n.. \} \times s ' \{n.. \}$ **for** p **unfolding** *dist-norm* **using** *that* **by** *measurable*

moreover have *countable* ($s \cdot \{n..\} \times s \cdot \{n..\}$) **by** (*intro countable-SIGMA*
countable-image, auto)
ultimately show *?thesis* **unfolding** * **by** (*auto intro!:* *borel-measurable-cSUP*
bdd)
qed

lemma *integrable-bound-diameter:*

fixes $f :: 'a \Rightarrow \text{real}$
assumes *integrable* M f
and [*measurable*]: $\bigwedge i. (s\ i) \in \text{borel-measurable } M$
and $\bigwedge x\ i. x \in \text{space } M \implies \text{norm } (s\ i\ x) \leq f\ x$
shows *integrable* M ($\lambda x. \text{diameter } \{s\ i\ x \mid i. n \leq i\}$)
proof –
have $\{s\ i\ x \mid i. N \leq i\} \neq \{\}$ **for** $x\ N$ **by** *blast*
hence *diameter-SUP*: $\text{diameter } \{s\ i\ x \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s\ i\ x) (s\ j\ x))$ **for** $x\ N$ **unfolding** *diameter-def* **by** (*auto intro!:* *arg-cong[of -*
- Sup])
{
fix x **assume** $x: x \in \text{space } M$
let $?S = (\lambda(i, j). \text{dist } (s\ i\ x) (s\ j\ x)) \cdot (\{n..\} \times \{n..\})$
have *case-prod dist* $\cdot (\{s\ i\ x \mid i. n \leq i\} \times \{s\ i\ x \mid i. n \leq i\}) = (\lambda(i, j). \text{dist } (s\ i\ x) (s\ j\ x)) \cdot (\{n..\} \times \{n..\})$ **by** *fast*
hence *: $\text{diameter } \{s\ i\ x \mid i. n \leq i\} = \text{Sup } ?S$ **using** *diameter-SUP* **by** (*simp*
add: case-prod-beta')

have *bounded* $?S$ **by** (*rule bounded-imp-dist-bounded[OF bounded[OF x]]*)
hence $\text{Sup } S\text{-nonneg:} 0 \leq \text{Sup } ?S$ **by** (*auto intro!:* *cSup-upper2 x bounded-imp-bdd-above*)

have $\text{dist } (s\ i\ x) (s\ j\ x) \leq 2 * f\ x$ **for** $i\ j$ **by** (*intro dist-triangle2[THEN*
order-trans, of - 0]) (*metis norm-conv-dist assms(3) x add-mono mult-2*)
hence $\forall c \in ?S. c \leq 2 * f\ x$ **by** *force*
hence $\text{Sup } ?S \leq 2 * f\ x$ **by** (*intro cSup-least, auto*)
hence $\text{norm } (\text{Sup } ?S) \leq 2 * \text{norm } (f\ x)$ **using** *Sup-S-nonneg* **by** *auto*
also have $\dots = \text{norm } (2 *_R f\ x)$ **by** *simp*
finally have $\text{norm } (\text{diameter } \{s\ i\ x \mid i. n \leq i\}) \leq \text{norm } (2 *_R f\ x)$ **unfolding**
 * .
}
hence $\text{AE } x \text{ in } M. \text{norm } (\text{diameter } \{s\ i\ x \mid i. n \leq i\}) \leq \text{norm } (2 *_R f\ x)$ **by** *blast*
thus *integrable* M ($\lambda x. \text{diameter } \{s\ i\ x \mid i. n \leq i\}$) **using** *borel-measurable-diameter*
by (*intro Bochner-Integration.integrable-bound[OF assms(1)[THEN integrable-scaleR-right[of*
2]]], measurable)
qed
end

end

theory *Set-Integral-Addendum*

imports *HOL-Analysis.Set-Integral Bochner-Integration-Addendum*
begin

lemma *set-integral-scaleR-left*:

assumes $A \in \text{sets } M \text{ } c \neq 0 \implies \text{integrable } M f$
shows $\text{LINT } t:A|M. f t *_{\mathbb{R}} c = (\text{LINT } t:A|M. f t) *_{\mathbb{R}} c$
unfolding *set-lebesgue-integral-def*
using *integrable-mult-indicator[OF assms]*
by (*subst integral-scaleR-left[symmetric], auto*)

lemma *nn-set-integral-eq-set-integral*:

assumes [*measurable*]:*integrable* $M f$
and $\forall x \in A \text{ in } M. 0 \leq f x \implies A \in \text{sets } M$
shows $(\int^+ x \in A. f x \partial M) = (\int x \in A. f x \partial M)$

proof –

have $(\int^+ x. \text{indicator } A x *_{\mathbb{R}} f x \partial M) = (\int x \in A. f x \partial M)$
unfolding *set-lebesgue-integral-def* **using** *assms(2)* **by** (*intro nn-integral-eq-integral[of*
*- \lambda x. \text{indicator } A x *_{\mathbb{R}} f x]*, *blast intro: assms integrable-mult-indicator, fastforce*)
moreover have $(\int^+ x. \text{indicator } A x *_{\mathbb{R}} f x \partial M) = (\int^+ x \in A. f x \partial M)$ **by** (*metis*
ennreal-0 indicator-simps(1) indicator-simps(2) mult.commute mult-1 mult-zero-left
real-scaleR-def)
ultimately show *?thesis* **by** *argov*
qed

lemma *set-integral-restrict-space*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{banach}, \text{second-countable-topology}\}$
assumes $\Omega \cap \text{space } M \in \text{sets } M$
shows $\text{set-lebesgue-integral } (\text{restrict-space } M \Omega) A f = \text{set-lebesgue-integral } M A$
 $(\lambda x. \text{indicator } \Omega x *_{\mathbb{R}} f x)$
unfolding *set-lebesgue-integral-def*
by (*subst integral-restrict-space, auto intro!: integrable-mult-indicator assms simp:*
mult.commute)

lemma *set-integral-const*:

fixes $c :: 'b :: \{\text{banach}, \text{second-countable-topology}\}$
assumes $A \in \text{sets } M \text{ } \text{emeasure } M A \neq \infty$
shows $\text{set-lebesgue-integral } M A (\lambda \cdot. c) = \text{measure } M A *_{\mathbb{R}} c$
unfolding *set-lebesgue-integral-def*
using *assms* **by** (*metis has-bochner-integral-indicator has-bochner-integral-integral-eq*
infinity-ennreal-def less-top)

lemma *set-integral-mono-banach*:

fixes $f g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}, \text{linorder-topology}, \text{ordered-real-vector}\}$
assumes *set-integrable* $M A f$ *set-integrable* $M A g$
 $\bigwedge x. x \in A \implies f x \leq g x$
shows $(\text{LINT } x:A|M. f x) \leq (\text{LINT } x:A|M. g x)$
using *assms* **unfolding** *set-integrable-def set-lebesgue-integral-def*
by (*auto intro: integral-mono-banach split: split-indicator*)

lemma *set-integral-mono-AE-banach*:

fixes $f g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}, \text{linorder-topology}, \text{ordered-real-vector}\}$

```

dered-real-vector}
  assumes set-integrable M A f set-integrable M A g AE x ∈ A in M. f x ≤ g x
  shows set-lebesgue-integral M A f ≤ set-lebesgue-integral M A g using assms
unfolding set-lebesgue-integral-def by (auto simp add: set-integrable-def intro!:
integral-mono-AE-banach[of M λx. indicator A x *_R f x λx. indicator A x *_R g x],
simp add: indicator-def)

end
theory Sigma-Finite-Measure-Addendum
imports Set-Integral-Addendum
begin

lemma balls-countable-basis:
  obtains D :: 'a :: {metric-space, second-countable-topology} set
  where topological-basis (case-prod ball ' (D × (ℚ ∩ {0<..})))
    and countable D
    and D ≠ {}
proof -
  obtain D :: 'a set where dense-subset: countable D D ≠ {} [[open U; U ≠ {}]]
  ⇒ ∃ y ∈ D. y ∈ U for U using countable-dense-exists by blast
  have topological-basis (case-prod ball ' (D × (ℚ ∩ {0<..})))
  proof (intro topological-basis-iff[THEN iffD2], fast, clarify)
    fix U and x :: 'a assume asm: open U x ∈ U
    obtain e where e: e > 0 ball x e ⊆ U using asm openE by blast
    obtain y where y: y ∈ D y ∈ ball x (e / 3) using dense-subset(3)[OF open-ball,
of x e / 3] centre-in-ball[THEN iffD2, OF divide-pos-pos[OF e(1), of 3]] by force
    obtain r where r: r ∈ ℚ ∩ {e/3<.. $e/2$ } unfolding Rats-def using of-rat-dense[OF
divide-strict-left-mono[OF - e(1)], of 2 3] by auto

    have *: x ∈ ball y r using r y by (simp add: dist-commute)
    hence ball y r ⊆ U using r by (intro order-trans[OF - e(2)], simp, metric)
    moreover have ball y r ∈ (case-prod ball ' (D × (ℚ ∩ {0<..}))) using y(1)
  r by force
    ultimately show ∃ B' ∈ (case-prod ball ' (D × (ℚ ∩ {0<..}))). x ∈ B' ∧ B' ⊆
U using * by meson
  qed
  thus ?thesis using that dense-subset by blast
qed

```

```

context sigma-finite-measure
begin

```

```

lemma sigma-finite-measure-induct[case-names finite-measure, consumes 0]:
  assumes ∧(N :: 'a measure) Ω. finite-measure N
    ⇒ N = restrict-space M Ω
    ⇒ Ω ∈ sets M
    ⇒ emeasure N Ω ≠ ∞
    ⇒ emeasure N Ω ≠ 0

```

$\implies \text{almost-everywhere } N \ Q$

and $[measurable]: \text{Measurable.pred } M \ Q$
shows $\text{almost-everywhere } M \ Q$
proof –
have $*$: $\text{almost-everywhere } N \ Q$ **if** $\text{finite-measure } N \ N = \text{restrict-space } M \ \Omega \ \Omega$
 $\in \text{sets } M \ \text{emeasure } N \ \Omega \neq \infty$ **for** $N \ \Omega$ **using** *that* **by** (*cases* $\text{emeasure } N \ \Omega = 0$,
auto *intro*: $\text{emeasure-0-AE assms}(1)$)

obtain $A :: \text{nat} \Rightarrow 'a \text{ set}$ **where** $A: \text{range } A \subseteq \text{sets } M \ (\bigcup i. A \ i) = \text{space } M$ **and**
 $\text{emeasure-finite: emeasure } M \ (A \ i) \neq \infty$ **for** i **using** $\text{sigma-finite by metis}$
note $A(1)[measurable]$
have $\text{space-restr: space } (\text{restrict-space } M \ (A \ i)) = A \ i$ **for** i **unfolding** $\text{space-restrict-space}$
by *simp*
 $\{$
 fix i
 have $*$: $\{x \in A \ i \cap \text{space } M. \ Q \ x\} = \{x \in \text{space } M. \ Q \ x\} \cap (A \ i)$ **by** *fast*
 have $\text{Measurable.pred } (\text{restrict-space } M \ (A \ i)) \ Q$ **using** A **by** (*intro measurableI*,
auto simp add: space-restr intro!: sets-restrict-space-iff[THEN iffD2], *measurable*,
auto)
 $\}$
note $\text{this}[measurable]$
 $\{$
 fix i
 have $\text{finite-measure } (\text{restrict-space } M \ (A \ i))$ **using** $\text{emeasure-finite by (intro$
finite-measureI, subst space-restr, subst emeasure-restrict-space, auto)
 hence $\text{emeasure } (\text{restrict-space } M \ (A \ i)) \ \{x \in A \ i. \neg Q \ x\} = 0$ **using** $\text{emeasure-finite by (intro AE-iff-measurable[THEN iffD1, OF - - *], measurable, subst$
space-restr[symmetric], intro sets.top, auto simp add: emeasure-restrict-space)
 hence $\text{emeasure } M \ \{x \in A \ i. \neg Q \ x\} = 0$ **by** (*subst emeasure-restrict-space[symmetric]*,
auto)
 $\}$
 hence $\text{emeasure } M \ (\bigcup i. \{x \in A \ i. \neg Q \ x\}) = 0$ **by** (*intro emeasure-UN-eq-0, auto*)
 moreover $\text{have } (\bigcup i. \{x \in A \ i. \neg Q \ x\}) = \{x \in \text{space } M. \neg Q \ x\}$ **using** A **by**
auto
 ultimately show *?thesis* **by** (*intro AE-iff-measurable[THEN iffD2]*, *auto*)
qed

lemma *averaging-theorem*:

fixes $f :: \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes $[measurable]: \text{integrable } M \ f$
and $\text{closed: closed } S$
and $\bigwedge A. A \in \text{sets } M \implies \text{measure } M \ A > 0 \implies (1 / \text{measure } M \ A) *_{\mathbb{R}}$
 $\text{set-lebesgue-integral } M \ A \ f \in S$
shows $\text{AE } x \text{ in } M. f \ x \in S$
proof (*induct rule: sigma-finite-measure-induct*)
case (*finite-measure* $N \ \Omega$)

interpret *finite-measure* N **by** (*rule finite-measure*)

have *integrable[measurable]*: *integrable* $N f$ **using** *assms finite-measure* **by** (*auto simp: integrable-restrict-space integrable-mult-indicator*)

have *average*: $(1 / \text{Sigma-Algebra.measure } N A) *_{\mathbb{R}} \text{set-lebesgue-integral } N A f \in S$ **if** $A \in \text{sets } N$ *measure* $N A > 0$ **for** A

proof –

have $*$: $A \in \text{sets } M$ **using** *that* **by** (*simp add: sets-restrict-space-iff finite-measure*)

have $A = A \cap \Omega$ **by** (*metis finite-measure(2,3) inf.orderE sets.sets-into-space space-restrict-space that(1)*)

hence *set-lebesgue-integral* $N A f = \text{set-lebesgue-integral } M A f$ **unfolding** *finite-measure* **by** (*subst set-integral-restrict-space, auto simp add: finite-measure set-lebesgue-integral-def indicator-inter-arith[symmetric]*)

moreover **have** *measure* $N A = \text{measure } M A$ **using** *that* **by** (*auto intro!: measure-restrict-space simp add: finite-measure sets-restrict-space-iff*)

ultimately show *?thesis* **using** *that* $*$ *assms(3)* **by** *presburger*

qed

obtain $D :: 'b \text{ set}$ **where** *balls-basis: topological-basis* (*case-prod ball ' (D × (Q ∩ {0<..}))*) **and** *countable-D: countable* D **using** *balls-countable-basis* **by** *blast*

have *countable-balls: countable* (*case-prod ball ' (D × (Q ∩ {0<..}))*) **using** *countable-rat countable-D* **by** *blast*

obtain B **where** *B-balls: B ⊆ case-prod ball ' (D × (Q ∩ {0<..})) ∪ B = −S* **using** *topological-basis[THEN iffD1, OF balls-basis] open-Compl[OF assms(2)]* **by** *meson*

hence *countable-B: countable* B **using** *countable-balls countable-subset* **by** *fast*

define b **where** $b = \text{from-nat-into } (B \cup \{\{\}\})$

have $B \cup \{\{\}\} \neq \{\}$ **by** *simp*

have *range-b: range* $b = B \cup \{\{\}\}$ **using** *countable-B* **by** (*auto simp add: b-def intro!: range-from-nat-into*)

have *open-b: open* $(b i)$ **for** i **unfolding** *b-def* **using** *B-balls open-ball from-nat-into[of B ∪ { { } } i]* **by** *force*

have *Union-range-b: ∪(range b) = −S* **using** *B-balls range-b* **by** *simp*

{

fix $v r$ **assume** *ball-in-Compl: ball* $v r \subseteq -S$

define A **where** $A = f^{-1} \text{ball } v r \cap \text{space } N$

have *dist-less: dist* $(f x) v < r$ **if** $x \in A$ **for** x **using** *that* **unfolding** *A-def vimage-def* **by** (*simp add: dist-commute*)

hence *AE-less: AE* $x \in A$ *in* N . *norm* $(f x - v) < r$ **by** (*auto simp add: dist-norm*)

have $*$: $A \in \text{sets } N$ **unfolding** *A-def* **by** *simp*

have *emeasure* $N A = 0$

proof –

{

assume *asm: emeasure* $N A > 0$

hence *measure-pos: measure* $N A > 0$ **unfolding** *emeasure-eq-measure* **by**

simp

hence $(1 / \text{measure } N \ A) *_{\mathcal{R}} \text{set-lebesgue-integral } N \ A \ f - v = (1 / \text{measure } N \ A) *_{\mathcal{R}} \text{set-lebesgue-integral } N \ A \ (\lambda x. f \ x - v)$ **using** *integrable integrable-const* * **by** $(\text{subst set-integral-diff } (2), \text{auto simp add: set-integrable-def set-integral-const}[OF *])$ *algebra-simps intro!: integrable-mult-indicator*

moreover have $\text{norm } (\int_{x \in A. (f \ x - v) \partial N}) \leq (\int_{x \in A. \text{norm } (f \ x - v) \partial N})$ **using** * **by** $(\text{auto intro!: integral-norm-bound}[of \ N \ \lambda x. \text{indicator } A \ x] *_{\mathcal{R}} (f \ x - v), \text{THEN order-trans})$ *integrable-mult-indicator integrable simp add: set-lebesgue-integral-def*

ultimately have $\text{norm } ((1 / \text{measure } N \ A) *_{\mathcal{R}} \text{set-lebesgue-integral } N \ A \ f - v) \leq \text{set-lebesgue-integral } N \ A \ (\lambda x. \text{norm } (f \ x - v)) / \text{measure } N \ A$ **using** *asm* **by** $(\text{auto intro: divide-right-mono})$

also have $\dots < \text{set-lebesgue-integral } N \ A \ (\lambda x. r) / \text{measure } N \ A$

unfolding *set-lebesgue-integral-def*

using *asm* * *integrable integrable-const AE-less measure-pos*

by $(\text{intro divide-strict-right-mono integral-less-AE}[of \ - \ A] \text{integrable-mult-indicator}) +$ $(\text{fastforce simp add: dist-less dist-norm indicator-def}) +$

also have $\dots = r$ **using** * *measure-pos* **by** $(\text{simp add: set-integral-const})$

finally have $\text{dist } ((1 / \text{measure } N \ A) *_{\mathcal{R}} \text{set-lebesgue-integral } N \ A \ f) \ v < r$ **by** (subst dist-norm)

hence *False* **using** *average[OF * measure-pos]* **by** $(\text{metis ComplD dist-commute in-mono mem-ball ball-in-Compl})$

}

thus *?thesis* **by** *fastforce*

qed

}

note * = *this*

{

fix *b'* **assume** *b' ∈ B*

hence *ball-subset-Compl: b' ⊆ -S* **and** *ball-radius-pos: ∃ v ∈ D. ∃ r > 0. b' = ball v r* **using** *B-balls* **by** $(\text{blast}, \text{fast})$

}

note ** = *this*

hence $\text{emeasure } N \ (f - ' b \ i \cap \text{space } N) = 0$ **for** *i* **by** $(\text{cases } b \ i = \{\}, \text{simp})$ $(\text{metis UnE singletonD } * \text{range-b}[THEN \text{eq-refl}, THEN \text{range-subsetD}])$

hence $\text{emeasure } N \ (\bigcup i. f - ' b \ i \cap \text{space } N) = 0$ **using** *open-b* **by** $(\text{intro emeasure-UN-eq-0}) \text{fastforce} +$

moreover have $(\bigcup i. f - ' b \ i \cap \text{space } N) = f - ' (\bigcup (\text{range } b)) \cap \text{space } N$ **by** *blast*

ultimately have $\text{emeasure } N \ (f - ' (-S) \cap \text{space } N) = 0$ **using** *Union-range-b* **by** *argo*

hence $\text{AE } x \text{ in } N. f \ x \notin -S$ **using** *open-Compl[OF assms(2)]* **by** $(\text{intro AE-iff-measurable}[THEN \text{iffD2}], \text{auto})$

thus *?case* **by** *force*

qed $(\text{simp add: pred-sets2}[OF \text{borel-closed}] \text{assms}(2))$

lemma *density-zero:*

fixes *f::'a ⇒ 'b::\{second-countable-topology, banach\}*

assumes *integrable M f*

and *density-0*: $\bigwedge A. A \in \text{sets } M \implies \text{set-lebesgue-integral } M A f = 0$
shows *AE* x in M . $f x = 0$
using *averaging-theorem*[*OF* *assms*(1), of {0}] *assms*(2)
by (*simp* add: *scaleR-nonneg-nonneg*)

lemma *density-unique*:
fixes $f f' :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes *integrable* $M f$ *integrable* $M f'$
and *density-eq*: $\bigwedge A. A \in \text{sets } M \implies \text{set-lebesgue-integral } M A f = \text{set-lebesgue-integral } M A f'$
shows *AE* x in M . $f x = f' x$
proof –
 {
fix A **assume** *asm*: $A \in \text{sets } M$
hence *LINT* $x | M$. *indicat-real* $A x *_R (f x - f' x) = 0$ **using** *density-eq*
assms(1,2) **by** (*simp* add: *set-lebesgue-integral-def algebra-simps Bochner-Integration.integral-diff*[*OF* *integrable-mult-indicator*(1,1)])
 }
thus ?thesis **using** *density-zero*[*OF* *Bochner-Integration.integrable-diff*[*OF* *assms*(1,2)]]
by (*simp* add: *set-lebesgue-integral-def*)
qed

lemma *density-nonneg*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes *integrable* $M f$
and $\bigwedge A. A \in \text{sets } M \implies \text{set-lebesgue-integral } M A f \geq 0$
shows *AE* x in M . $f x \geq 0$
using *averaging-theorem*[*OF* *assms*(1), of {0..}, *OF* *closed-atLeast*] *assms*(2)
by (*simp* add: *scaleR-nonneg-nonneg*)

corollary *integral-nonneg-AE-eq-0-iff-AE*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes $f[\text{measurable}]$: *integrable* $M f$ **and** *nonneg*: *AE* x in M . $0 \leq f x$
shows *integral*^L $M f = 0 \iff (\text{AE } x \text{ in } M. f x = 0)$
proof
assume *: *integral*^L $M f = 0$
 {
fix A **assume** *asm*: $A \in \text{sets } M$
have $0 \leq \text{integral}^L M (\lambda x. \text{indicator } A x *_R f x)$ **using** *nonneg* **by** (*subst integr-zero*[of M , *symmetric*], *intro integral-mono-AE-banach integrable-mult-indicator asm f integrable-zero*, *auto simp* add: *indicator-def*)
moreover **have** $\dots \leq \text{integral}^L M f$ **using** *nonneg* **by** (*intro integral-mono-AE-banach integrable-mult-indicator asm f*, *auto simp* add: *indicator-def*)
ultimately **have** *set-lebesgue-integral* $M A f = 0$ **unfolding** *set-lebesgue-integral-def* **using** * **by** *force*
 }
thus *AE* x in M . $f x = 0$ **by** (*intro density-zero f*, *blast*)
qed (*auto simp* add: *integral-eq-zero-AE*)

corollary *integral-eq-mono-AE-eq-AE*:

fixes $f\ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$

assumes $\text{integrable } M\ f\ \text{integrable } M\ g\ \text{integral}^L\ M\ f = \text{integral}^L\ M\ g\ \text{AE } x\ \text{in } M.\ f\ x \leq g\ x$

shows $\text{AE } x\ \text{in } M.\ f\ x = g\ x$

proof –

define h **where** $h = (\lambda x. g\ x - f\ x)$

have $\text{AE } x\ \text{in } M.\ h\ x = 0$ **unfolding** $h\text{-def}$ **using** assms **by** $(\text{subst integral-nonneg-AE-eq-0-iff-AE[symmetric]})\ \text{auto}$

then show $?thesis$ **unfolding** $h\text{-def}$ **by** auto

qed

end

end

theory *Filtered-Measure*

imports *HOL-Probability.Conditional-Expectation*

begin

0.5 Filtered Measure

locale *filtered-measure* =

fixes $M\ F$ **and** $t_0 :: 'b :: \{\text{second-countable-topology, linorder-topology}\}$

assumes $\text{subalgebra: } \bigwedge i. t_0 \leq i \implies \text{subalgebra } M\ (F\ i)$

and $\text{sets-}F\text{-mono: } \bigwedge i\ j. t_0 \leq i \implies i \leq j \implies \text{sets } (F\ i) \leq \text{sets } (F\ j)$

begin

lemma *space-F*:

assumes $t_0 \leq i$

shows $\text{space } (F\ i) = \text{space } M$

using *subalgebra* assms **by** $(\text{simp add: subalgebra-def})$

lemma *subalgebra-F*:

assumes $t_0 \leq i\ i \leq j$

shows $\text{subalgebra } (F\ j)\ (F\ i)$

unfolding *subalgebra-def* **using** assms **by** $(\text{simp add: space-F sets-F-mono})$

lemma *borel-measurable-mono*:

assumes $t_0 \leq i\ i \leq j$

shows $\text{borel-measurable } (F\ i) \subseteq \text{borel-measurable } (F\ j)$

unfolding *subset-iff* **by** $(\text{metis assms subalgebra-F measurable-from-subalg})$

end

0.6 Filtered Sigma Finite Measure

The locale presented here is a generalization of the *sigma-finite-subalgebra* for a particular filtration.

locale *sigma-finite-filtered-measure* = *filtered-measure* +
assumes *sigma-finite*: *sigma-finite-subalgebra* *M* (*F* *t*₀)

lemma (**in** *sigma-finite-filtered-measure*) *sigma-finite-subalgebra-F*[*intro*]:
assumes *t*₀ ≤ *i*
shows *sigma-finite-subalgebra* *M* (*F* *i*)
using *assms* **by** (*metis dual-order.refl sets-F-mono sigma-finite sigma-finite-subalgebra.nested-subalg-is-sigma-subalgebra subalgebra-def*)

0.6.1 Typed locales

locale *nat-filtered-measure* = *filtered-measure* *M* *F* 0 **for** *M* **and** *F* :: *nat* ⇒ -
locale *real-filtered-measure* = *filtered-measure* *M* *F* 0 **for** *M* **and** *F* :: *real* ⇒ -

context *nat-filtered-measure*
begin

lemma *space-F*: *space* (*F* *i*) = *space* *M*
using *subalgebra* **by** (*simp add: subalgebra-def*)

lemma *subalgebra-F*:
assumes *i* ≤ *j*
shows *subalgebra* (*F* *j*) (*F* *i*)
unfolding *subalgebra-def* **using** *assms* **by** (*simp add: space-F sets-F-mono*)

lemma *borel-measurable-mono*:
assumes *i* ≤ *j*
shows *borel-measurable* (*F* *i*) ⊆ *borel-measurable* (*F* *j*)
unfolding *subset-iff* **by** (*metis assms subalgebra-F measurable-from-subalg*)

end

locale *nat-sigma-finite-filtered-measure* = *sigma-finite-filtered-measure* *M* *F* 0 **for** *M* **and** *F* :: *nat* ⇒ -
locale *real-sigma-finite-filtered-measure* = *sigma-finite-filtered-measure* *M* *F* 0 **for** *M* **and** *F* :: *real* ⇒ -

sublocale *nat-sigma-finite-filtered-measure* ⊆ *sigma-finite-subalgebra* *M* *F* *i* **by** *blast*

0.6.2 Constant Filtration

lemma *filtered-measure-constant-filtration*:
assumes *subalgebra* *M* *F*
shows *filtered-measure* *M* (λ-. *F*) *t*₀

```

using assms by (unfold-locales) (auto simp add: subalgebra-def)

sublocale sigma-finite-subalgebra  $\subseteq$  constant-filtration: sigma-finite-filtered-measure
M  $\lambda$ - :: 't :: {second-countable-topology, linorder-topology}. F t0
  using subalg by (unfold-locales) (auto simp add: subalgebra-def)

end
theory Conditional-Expectation-Banach
imports HOL-Probability.Conditional-Expectation Sigma-Finite-Measure-Addendum
begin

definition has-cond-exp :: 'a measure  $\Rightarrow$  'a measure  $\Rightarrow$  ('a  $\Rightarrow$  'b)  $\Rightarrow$  ('a  $\Rightarrow$  'b::{real-normed-vector,
second-countable-topology})  $\Rightarrow$  bool where
  has-cond-exp M F f g = (( $\forall A \in \text{sets } F. (\int x \in A. f x \partial M) = (\int x \in A. g x \partial M)$ )
     $\wedge$  integrable M f
     $\wedge$  integrable M g
     $\wedge$  g  $\in$  borel-measurable F)

lemma has-cond-expI':
  assumes  $\bigwedge A. A \in \text{sets } F \implies (\int x \in A. f x \partial M) = (\int x \in A. g x \partial M)$ 
    integrable M f
    integrable M g
    g  $\in$  borel-measurable F
  shows has-cond-exp M F f g
  using assms unfolding has-cond-exp-def by simp

lemma has-cond-expD:
  assumes has-cond-exp M F f g
  shows  $\bigwedge A. A \in \text{sets } F \implies (\int x \in A. f x \partial M) = (\int x \in A. g x \partial M)$ 
    integrable M f
    integrable M g
    g  $\in$  borel-measurable F
  using assms unfolding has-cond-exp-def by simp+

definition cond-exp :: 'a measure  $\Rightarrow$  'a measure  $\Rightarrow$  ('a  $\Rightarrow$  'b)  $\Rightarrow$  ('a  $\Rightarrow$  'b::{banach,
second-countable-topology}) where
  cond-exp M F f = (if  $\exists g. \text{has-cond-exp } M \text{ } F \text{ } f \text{ } g$  then (SOME g. has-cond-exp M
F f g) else ( $\lambda$ -. 0))

lemma borel-measurable-cond-exp[measurable]: cond-exp M F f  $\in$  borel-measurable
F
  by (metis cond-exp-def someI has-cond-exp-def borel-measurable-const)

lemma integrable-cond-exp[intro]: integrable M (cond-exp M F f)
  by (metis cond-exp-def has-cond-expD(3) integrable-zero someI)

```

lemma *set-integrable-cond-exp*[intro]:
assumes $A \in \text{sets } M$
shows *set-integrable* $M A$ (*cond-exp* $M F f$) **using** *integrable-mult-indicator*[*OF* *assms integrable-cond-exp*, *of F f*] **by** (*auto simp add: set-integrable-def intro!: integrable-mult-indicator*[*OF assms integrable-cond-exp*])

context *sigma-finite-subalgebra*
begin

lemma *borel-measurable-cond-exp'*[*measurable*]: *cond-exp* $M F f \in \text{borel-measurable } M$
by (*metis cond-exp-def someI has-cond-exp-def borel-measurable-const subalg measurable-from-subalg*)

lemma *cond-exp-null*:
assumes $\nexists g. \text{has-cond-exp } M F f g$
shows *cond-exp* $M F f = (\lambda-. 0)$
unfolding *cond-exp-def* **using** *assms* **by** *argo*

lemma *has-cond-exp-nested-subalg*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes *subalgebra* $G F$ *has-cond-exp* $M F f h$ *has-cond-exp* $M G f h'$
shows *has-cond-exp* $M F h' h$
by (*intro has-cond-expI'*) (*metis assms has-cond-expD in-mono subalgebra-def*) +

lemma *has-cond-exp-charact*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes *has-cond-exp* $M F f g$
shows *has-cond-exp* $M F f$ (*cond-exp* $M F f$)
 $\text{AE } x \text{ in } M. \text{cond-exp } M F f x = g x$

proof –
show *cond-exp*: *has-cond-exp* $M F f$ (*cond-exp* $M F f$) **using** *assms someI cond-exp-def* **by** *metis*
let $?MF = \text{restr-to-subalg } M F$
interpret *sigma-finite-measure* $?MF$ **by** (*rule sigma-fin-subalg*)
{
fix A **assume** $A \in \text{sets } ?MF$
then have [*measurable*]: $A \in \text{sets } F$ **using** *sets-restr-to-subalg*[*OF subalg*] **by** *simp*
have $(\int x \in A. g x \partial ?MF) = (\int x \in A. g x \partial M)$ **using** *assms subalg* **by** (*auto simp add: integral-subalgebra2 set-lebesgue-integral-def dest!: has-cond-expD*)
also have $\dots = (\int x \in A. \text{cond-exp } M F f x \partial M)$ **using** *assms cond-exp* **by** (*simp add: has-cond-exp-def*)
also have $\dots = (\int x \in A. \text{cond-exp } M F f x \partial ?MF)$ **using** *subalg* **by** (*auto simp add: integral-subalgebra2 set-lebesgue-integral-def*)
finally have $(\int x \in A. g x \partial ?MF) = (\int x \in A. \text{cond-exp } M F f x \partial ?MF)$ **by** *simp*
}
hence $\text{AE } x \text{ in } ?MF. \text{cond-exp } M F f x = g x$ **using** *cond-exp assms subalg* **by**

(intro density-unique, auto dest: has-cond-expD intro!: integrable-in-subalg)
 then show $\text{AE } x \text{ in } M. \text{ cond-exp } M F f x = g x$ using $\text{AE-restr-to-subalg}[OF \text{ subalg}]$ by simp
 qed

lemma cond-exp-charact:
 fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
 assumes $\bigwedge A. A \in \text{sets } F \implies (\int x \in A. f x \partial M) = (\int x \in A. g x \partial M)$
 integrable $M f$
 integrable $M g$
 $g \in \text{borel-measurable } F$
 shows $\text{AE } x \text{ in } M. \text{ cond-exp } M F f x = g x$
 by (intro has-cond-exp-charact has-cond-expI' assms) auto

corollary cond-exp-F-meas[intro, simp]:
 fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
 assumes integrable $M f$
 $f \in \text{borel-measurable } F$
 shows $\text{AE } x \text{ in } M. \text{ cond-exp } M F f x = f x$
 by (rule cond-exp-charact, auto intro: assms)

Congruence

lemma has-cond-exp-cong:
 assumes integrable $M f \bigwedge x. x \in \text{space } M \implies f x = g x$ has-cond-exp $M F g h$
 shows has-cond-exp $M F f h$
proof (intro has-cond-expI'[OF - assms(1)], goal-cases)
 case (1 A)
 hence $\text{set-lebesgue-integral } M A f = \text{set-lebesgue-integral } M A g$ by (intro set-lebesgue-integral-cong)
 (meson assms(2) subalg in-mono subalgebra-def sets.sets-into-space subalgebra-def subsetD)+
 then show ?case using 1 assms(3) by (simp add: has-cond-exp-def)
qed (auto simp add: has-cond-expD[OF assms(3)])

lemma cond-exp-cong:
 fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
 assumes integrable $M f$ integrable $M g \bigwedge x. x \in \text{space } M \implies f x = g x$
 shows $\text{AE } x \text{ in } M. \text{ cond-exp } M F f x = \text{cond-exp } M F g x$
proof (cases $\exists h. \text{has-cond-exp } M F f h$)
 case True
 then obtain h where $h: \text{has-cond-exp } M F f h \text{ has-cond-exp } M F g h$ using
 has-cond-exp-cong assms by metis
 show ?thesis using $h[\text{THEN has-cond-exp-charact}(2)]$ by fastforce
 next
 case False
 moreover have $\nexists h. \text{has-cond-exp } M F g h$ using False has-cond-exp-cong assms
 by auto
 ultimately show ?thesis unfolding cond-exp-def by auto
qed

lemma *has-cond-exp-cong-AE*:
assumes *integrable M f AE x in M. f x = g x has-cond-exp M F g h*
shows *has-cond-exp M F f h*
using *assms(1,2) subalg subalgebra-def subset-iff*
by (*intro has-cond-expI'*, *subst set-lebesgue-integral-cong-AE[OF - assms(1)[THEN borel-measurable-integrable] borel-measurable-integrable(1)[OF has-cond-expD(2)[OF assms(3)]]]*)
(fast intro: has-cond-expD[OF assms(3)] integrable-cong-AE-imp[OF - - AE-symmetric])+

lemma *has-cond-exp-cong-AE'*:
assumes *h ∈ borel-measurable F AE x in M. h x = h' x has-cond-exp M F f h'*
shows *has-cond-exp M F f h*
using *assms(1, 2) subalg subalgebra-def subset-iff*
using *AE-restr-to-subalg2[OF subalg assms(2)] measurable-from-subalg*
by (*intro has-cond-expI'*, *subst set-lebesgue-integral-cong-AE[OF - measurable-from-subalg(1,1)[OF subalg], OF - assms(1) has-cond-expD(4)[OF assms(3)]]]*)
(fast intro: has-cond-expD[OF assms(3)] integrable-cong-AE-imp[OF - - AE-symmetric])+

lemma *cond-exp-cong-AE*:
fixes *f :: 'a ⇒ 'b::{second-countable-topology,banach}*
assumes *integrable M f integrable M g AE x in M. f x = g x*
shows *AE x in M. cond-exp M F f x = cond-exp M F g x*
proof (*cases ∃ h. has-cond-exp M F f h*)
case *True*
then obtain h where *h: has-cond-exp M F f h has-cond-exp M F g h* **using** *has-cond-exp-cong-AE assms* **by** (*metis (mono-tags, lifting) eventually-mono*)
show *?thesis* **using** *h[THEN has-cond-exp-charact(2)]* **by** *fastforce*
next
case *False*
moreover have *¬ h. has-cond-exp M F g h* **using** *False has-cond-exp-cong-AE assms* **by** *auto*
ultimately show *?thesis* **unfolding** *cond-exp-def* **by** *auto*
qed

lemma *has-cond-exp-real*:
fixes *f :: 'a ⇒ real*
assumes *integrable M f*
shows *has-cond-exp M F f (real-cond-exp M F f)*
by (*intro has-cond-expI'*, *auto intro!: real-cond-exp-intA assms*)

lemma *cond-exp-real[intro]*:
fixes *f :: 'a ⇒ real*
assumes *integrable M f*
shows *AE x in M. cond-exp M F f x = real-cond-exp M F f x*
using *has-cond-exp-charact has-cond-exp-real assms* **by** *blast*

lemma *cond-exp-cmult*:
fixes *f :: 'a ⇒ real*
assumes *integrable M f*

shows $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ (\lambda x.\ c * f\ x)\ x = c * cond\text{-}exp\ M\ F\ f\ x$
using $real\text{-}cond\text{-}exp\text{-}cmult[OF\ assms(1),\ of\ c]\ assms(1)[THEN\ cond\text{-}exp\text{-}real]$
 $assms(1)[THEN\ integrable\text{-}mult\text{-}right,\ THEN\ cond\text{-}exp\text{-}real,\ of\ c]$ **by** *fastforce*

Indicator functions

lemma *has-cond-exp-indicator*:

assumes $A \in sets\ M\ emeasure\ M\ A < \infty$
shows $has\text{-}cond\text{-}exp\ M\ F\ (\lambda x.\ indicat\text{-}real\ A\ x *_{\mathbb{R}} y)\ (\lambda x.\ real\text{-}cond\text{-}exp\ M\ F\ (indicator\ A)\ x *_{\mathbb{R}} y)$
proof (*intro has-cond-expI'*, *goal-cases*)
case (1 *B*)
have $\int_{x \in B} (indicat\text{-}real\ A\ x *_{\mathbb{R}} y)\ \partial M = (\int_{x \in B} indicat\text{-}real\ A\ x\ \partial M) *_{\mathbb{R}} y$ **using** *assms* **by** (*intro set-integral-scaleR-left, meson 1 in-mono subalg subalgebra-def, blast*)
also have $\dots = (\int_{x \in B} real\text{-}cond\text{-}exp\ M\ F\ (indicator\ A)\ x\ \partial M) *_{\mathbb{R}} y$ **using** 1 *assms* **by** (*subst real-cond-exp-intA, auto*)
also have $\dots = \int_{x \in B} (real\text{-}cond\text{-}exp\ M\ F\ (indicator\ A)\ x *_{\mathbb{R}} y)\ \partial M$ **using** *assms* **by** (*intro set-integral-scaleR-left[symmetric], meson 1 in-mono subalg subalgebra-def, blast*)
finally show ?*case* .
next
case 2
then show ?*case* **using** *integrable-scaleR-left integrable-real-indicator assms* **by** *blast*
next
case 3
show ?*case* **using** *assms* **by** (*intro integrable-scaleR-left, intro real-cond-exp-int, blast+*)
next
case 4
then show ?*case* **by** (*intro borel-measurable-scaleR, intro Conditional-Expectation.borel-measurable-cond-exp, simp*)
qed

lemma *cond-exp-indicator[intro]*:

fixes $y :: 'b :: \{second\text{-}countable\text{-}topology, banach\}$
assumes [*measurable*]: $A \in sets\ M\ emeasure\ M\ A < \infty$
shows $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ (\lambda x.\ indicat\text{-}real\ A\ x *_{\mathbb{R}} y)\ x = cond\text{-}exp\ M\ F\ (indicator\ A)\ x *_{\mathbb{R}} y$
proof –
have $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ (\lambda x.\ indicat\text{-}real\ A\ x *_{\mathbb{R}} y)\ x = real\text{-}cond\text{-}exp\ M\ F\ (indicator\ A)\ x *_{\mathbb{R}} y$ **using** *has-cond-exp-indicator[OF assms] has-cond-exp-charact* **by** *blast*
thus ?*thesis* **using** *cond-exp-real[OF integrable-real-indicator, OF assms]* **by** *fast-force*
qed

Addition

lemma *has-cond-exp-add*:

```

fixes  $f\ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
assumes  $\text{has-cond-exp } M\ F\ f\ f'\ \text{has-cond-exp } M\ F\ g\ g'$ 
shows  $\text{has-cond-exp } M\ F\ (\lambda x. f\ x + g\ x)\ (\lambda x. f'\ x + g'\ x)$ 
proof (intro  $\text{has-cond-expI'}$ , goal-cases)
  case (1  $A$ )
    have  $\int_{x \in A}. (f\ x + g\ x) \partial M = (\int_{x \in A}. f\ x\ \partial M) + (\int_{x \in A}. g\ x\ \partial M)$  using
       $\text{assms}[\text{THEN } \text{has-cond-expD}(2)]$  subalg 1 by (intro  $\text{set-integral-add}(2)$ , auto simp
       $\text{add: subalgebra-def set-integrable-def intro: integrable-mult-indicator}$ )
    also have  $\dots = (\int_{x \in A}. f'\ x\ \partial M) + (\int_{x \in A}. g'\ x\ \partial M)$  using  $\text{assms}[\text{THEN } \text{has-cond-expD}(1)[\text{OF } - 1]]$  by argo
    also have  $\dots = \int_{x \in A}. (f'\ x + g'\ x) \partial M$  using  $\text{assms}[\text{THEN } \text{has-cond-expD}(3)]$ 
      subalg 1 by (intro  $\text{set-integral-add}(2)[\text{symmetric}]$ , auto simp  $\text{add: subalgebra-def}$ 
       $\text{set-integrable-def intro: integrable-mult-indicator}$ )
    finally show ?case .
  next
    case 2
    then show ?case by (metis Bochner-Integration.integrable-add  $\text{assms has-cond-expD}(2)$ )
  next
    case 3
    then show ?case by (metis Bochner-Integration.integrable-add  $\text{assms has-cond-expD}(3)$ )
  next
    case 4
    then show ?case using  $\text{assms borel-measurable-add has-cond-expD}(4)$  by blast
qed

```

lemma *has-cond-exp-scaleR-right*:

```

fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
assumes  $\text{has-cond-exp } M\ F\ f\ f'$ 
shows  $\text{has-cond-exp } M\ F\ (\lambda x. c *_{\mathbb{R}} f\ x)\ (\lambda x. c *_{\mathbb{R}} f'\ x)$ 
using  $\text{has-cond-expD}[\text{OF } \text{assms}]$  by (intro  $\text{has-cond-expI'}$ , auto)

```

lemma *cond-exp-scaleR-right*:

```

fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
assumes  $\text{integrable } M\ f$ 
shows  $\text{AE } x \text{ in } M. \text{cond-exp } M\ F\ (\lambda x. c *_{\mathbb{R}} f\ x)\ x = c *_{\mathbb{R}} \text{cond-exp } M\ F\ f\ x$ 
proof (cases  $\exists f'. \text{has-cond-exp } M\ F\ f\ f'$ )
  case True
    then show ?thesis using  $\text{assms has-cond-exp-charact has-cond-exp-scaleR-right}$ 
by metis
  next
    case False
    show ?thesis
    proof (cases  $c = 0$ )
      case True
        then show ?thesis by simp
      next
        case  $c \neq 0$ : False
        have  $\nexists f'. \text{has-cond-exp } M\ F\ (\lambda x. c *_{\mathbb{R}} f\ x)\ f'$ 
        proof (standard, goal-cases)

```

```

    case 1
    then obtain f' where f': has-cond-exp M F (λx. c *R f x) f' by blast
    have has-cond-exp M F f (λx. inverse c *R f' x) using has-cond-expD[OF
f'] divideR-right[OF c-nonzero] assms by (intro has-cond-expI', auto)
    then show ?case using False by blast
  qed
  then show ?thesis using cond-exp-null[OF False] cond-exp-null by force
qed
qed

```

lemma *cond-exp-uminus*:

```

  fixes f :: 'a ⇒ 'b::{second-countable-topology,banach}
  assumes integrable M f
  shows AE x in M. cond-exp M F (λx. - f x) x = - cond-exp M F f x
  using cond-exp-scaleR-right[OF assms, of -1] by force

```

corollary *has-cond-exp-simple*:

```

  fixes f :: 'a ⇒ 'b::{second-countable-topology,banach}
  assumes simple-function M f emeasure M {y ∈ space M. f y ≠ 0} ≠ ∞
  shows has-cond-exp M F f (cond-exp M F f)
  using assms
proof (induction rule: integrable-simple-function-induct)
  case (cong f g)
  then show ?case using has-cond-exp-cong by (metis (no-types, opaque-lifting)
Bochner-Integration.integrable-cong has-cond-expD(2) has-cond-exp-charact(1))
next
  case (indicator A y)
  then show ?case using has-cond-exp-charact[OF has-cond-exp-indicator] by fast
next
  case (add u v)
  then show ?case using has-cond-exp-add has-cond-exp-charact(1) by blast
qed

```

lemma *cond-exp-contraction-real*:

```

  fixes f :: 'a ⇒ real
  assumes integrable[measurable]: integrable M f
  shows AE x in M. norm (cond-exp M F f x) ≤ cond-exp M F (λx. norm (f x)) x
proof-
  have int: integrable M (λx. norm (f x)) using assms by blast
  have *: AE x in M. 0 ≤ cond-exp M F (λx. norm (f x)) x using cond-exp-real[THEN
AE-symmetric, OF integrable-norm[OF integrable]] real-cond-exp-ge-c[OF integrable-norm[OF
integrable], of 0] norm-ge-zero by fastforce
  have **: A ∈ sets F ⇒ ∫ x∈A. |f x| ∂M = ∫ x∈A. real-cond-exp M F (λx.
norm (f x)) x ∂M for A unfolding real-norm-def using assms integrable-abs
real-cond-exp-intA by blast

```

```

  have norm-int: A ∈ sets F ⇒ (∫ x∈A. |f x| ∂M) = (∫+ x∈A. |f x| ∂M) for A
  using assms by (intro nn-set-integral-eq-set-integral[symmetric], blast, fastforce)
  (meson subalg subalgebra-def subsetD)

```


have $AE\ x\ in\ M. \text{real-cond-exp } M\ F\ (\lambda x. \text{norm } (f\ x))\ x \geq 0$ **using** *int real-cond-exp-ge-c*
by *force*

hence *cond-exp-norm-int*: $A \in \text{sets } F \implies (\int x \in A. \text{real-cond-exp } M\ F\ (\lambda x. \text{norm } (f\ x))\ x\ \partial M) = (\int x \in A. \text{real-cond-exp } M\ F\ (\lambda x. \text{norm } (f\ x))\ x\ \partial M)$ **for** A **using**
assms **by** (*intro nn-set-integral-eq-set-integral[symmetric]*, *blast*, *fastforce*) (*meson*
subalg subalgebra-def subsetD)

have $A \in \text{sets } F \implies \int x \in A. |f\ x| \partial M = \int x \in A. \text{real-cond-exp } M\ F\ (\lambda x. \text{norm } (f\ x))\ x\ \partial M$ **for** A **using** *** norm-int cond-exp-norm-int* **by** (*auto simp*
add: nn-integral-set-ennreal)

moreover **have** $(\lambda x. \text{ennreal } |f\ x|) \in \text{borel-measurable } M$ **by** *measurable*

moreover **have** $(\lambda x. \text{ennreal } (\text{real-cond-exp } M\ F\ (\lambda x. \text{norm } (f\ x))\ x)) \in \text{borel-measurable } F$ **by** *measurable*

ultimately **have** $AE\ x\ in\ M. \text{nn-cond-exp } M\ F\ (\lambda x. \text{ennreal } |f\ x|)\ x = \text{real-cond-exp } M\ F\ (\lambda x. \text{norm } (f\ x))\ x$ **by** (*intro nn-cond-exp-charact[THEN AE-symmetric]*,
auto)

hence $AE\ x\ in\ M. \text{nn-cond-exp } M\ F\ (\lambda x. \text{ennreal } |f\ x|)\ x \leq \text{cond-exp } M\ F\ (\lambda x. \text{norm } (f\ x))\ x$ **using** *cond-exp-real[OF int]* **by** *force*

moreover **have** $AE\ x\ in\ M. |\text{real-cond-exp } M\ F\ f\ x| = \text{norm } (\text{cond-exp } M\ F\ f\ x)$

unfolding *real-norm-def* **using** *cond-exp-real[OF assms]* *** by** *force*

ultimately **have** $AE\ x\ in\ M. \text{ennreal } (\text{norm } (\text{cond-exp } M\ F\ f\ x)) \leq \text{cond-exp } M\ F\ (\lambda x. \text{norm } (f\ x))\ x$ **using** *real-cond-exp-abs[OF assms[THEN borel-measurable-integrable]]*
by *fastforce*

hence $AE\ x\ in\ M. \text{enn2real } (\text{ennreal } (\text{norm } (\text{cond-exp } M\ F\ f\ x))) \leq \text{enn2real } (\text{cond-exp } M\ F\ (\lambda x. \text{norm } (f\ x))\ x)$ **using** *ennreal-le-iff2* **by** *force*

thus *?thesis* **using** *** **by** *fastforce*

qed

lemma *cond-exp-contraction-simple*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$

assumes *simple-function* $M\ f\ \text{emeasure } M\ \{y \in \text{space } M. f\ y \neq 0\} \neq \infty$

shows $AE\ x\ in\ M. \text{norm } (\text{cond-exp } M\ F\ f\ x) \leq \text{cond-exp } M\ F\ (\lambda x. \text{norm } (f\ x))\ x$

using *assms*

proof (*induction rule: integrable-simple-function-induct*)

case (*cong f g*)

hence *ae*: $AE\ x\ in\ M. f\ x = g\ x$ **by** *blast*

hence $AE\ x\ in\ M. \text{cond-exp } M\ F\ f\ x = \text{cond-exp } M\ F\ g\ x$ **using** *cong has-cond-exp-simple*
by (*subst cond-exp-cong-AE*) (*auto intro!: has-cond-expD(2)*)

hence $AE\ x\ in\ M. \text{norm } (\text{cond-exp } M\ F\ f\ x) = \text{norm } (\text{cond-exp } M\ F\ g\ x)$ **by**
force

moreover **have** $AE\ x\ in\ M. \text{cond-exp } M\ F\ (\lambda x. \text{norm } (f\ x))\ x = \text{cond-exp } M\ F\ (\lambda x. \text{norm } (g\ x))\ x$ **using** *ae cong has-cond-exp-simple* **by** (*subst cond-exp-cong-AE*)
(auto dest: has-cond-expD)

ultimately **show** *?case* **using** *cong(6)* **by** *fastforce*

next

case (*indicator A y*)

hence $AE\ x\ in\ M. \text{cond-exp } M\ F\ (\lambda a. \text{indicator } A\ a\ *_R\ y)\ x = \text{cond-exp } M\ F\ (\text{indicator } A)\ x\ *_R\ y$ **by** *blast*

hence *: $AE\ x\ in\ M.\ norm\ (cond\text{-}exp\ M\ F\ (\lambda a.\ indicat\text{-}real\ A\ a\ *_R\ y)\ x) \leq norm\ y$
 $*\ cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (indicat\text{-}real\ A\ x))\ x$ **using** $cond\text{-}exp\text{-}contraction\text{-}real[OF\ integrable\text{-}real\text{-}indicator,\ OF\ indicator]$ **by** *fastforce*

have $AE\ x\ in\ M.\ norm\ y\ *\ cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (indicat\text{-}real\ A\ x))\ x = norm\ y\ *\ real\text{-}cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (indicat\text{-}real\ A\ x))\ x$ **using** $cond\text{-}exp\text{-}real[OF\ integrable\text{-}real\text{-}indicator,\ OF\ indicator]$ **by** *fastforce*

moreover **have** $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ y\ *\ norm\ (indicat\text{-}real\ A\ x))\ x = real\text{-}cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ y\ *\ norm\ (indicat\text{-}real\ A\ x))\ x$ **using** *indicator* **by** (*intro cond-exp-real, auto*)

ultimately **have** $AE\ x\ in\ M.\ norm\ y\ *\ cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (indicat\text{-}real\ A\ x))\ x = cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ y\ *\ norm\ (indicat\text{-}real\ A\ x))\ x$ **using** $real\text{-}cond\text{-}exp\text{-}cmult[of\ \lambda x.\ norm\ (indicat\text{-}real\ A\ x)\ norm\ y]\ indicator$ **by** *fastforce*

moreover **have** $(\lambda x.\ norm\ y\ *\ norm\ (indicat\text{-}real\ A\ x)) = (\lambda x.\ norm\ (indicat\text{-}real\ A\ x\ *_R\ y))$ **by** *force*

ultimately **show** *?case* **using** * **by** *force*

next

case (*add u v*)

have $AE\ x\ in\ M.\ norm\ (cond\text{-}exp\ M\ F\ (\lambda a.\ u\ a + v\ a)\ x) = norm\ (cond\text{-}exp\ M\ F\ u\ x + cond\text{-}exp\ M\ F\ v\ x)$ **using** $has\text{-}cond\text{-}exp\text{-}charact(2)[OF\ has\text{-}cond\text{-}exp\text{-}add,\ OF\ has\text{-}cond\text{-}exp\text{-}simple(1,1),\ OF\ add(1,2,3,4)]$ **by** *fastforce*

moreover **have** $AE\ x\ in\ M.\ norm\ (cond\text{-}exp\ M\ F\ u\ x + cond\text{-}exp\ M\ F\ v\ x) \leq norm\ (cond\text{-}exp\ M\ F\ u\ x) + norm\ (cond\text{-}exp\ M\ F\ v\ x)$ **using** $norm\text{-}triangle\text{-}ineq$ **by** *blast*

moreover **have** $AE\ x\ in\ M.\ norm\ (cond\text{-}exp\ M\ F\ u\ x) + norm\ (cond\text{-}exp\ M\ F\ v\ x) \leq cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (u\ x))\ x + cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (v\ x))\ x$ **using** $add(6,7)$ **by** *fastforce*

moreover **have** $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (u\ x))\ x + cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (v\ x))\ x = cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (u\ x) + norm\ (v\ x))\ x$ **using** $integrable\text{-}simple\text{-}function[OF\ add(1,2)]\ integrable\text{-}simple\text{-}function[OF\ add(3,4)]$ **by** (*intro has-cond-exp-charact(2)[OF has-cond-exp-add[OF has-cond-exp-charact(1,1)], THEN AE-symmetric], auto intro: has-cond-exp-real*)

moreover **have** $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (u\ x) + norm\ (v\ x))\ x = cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (u\ x + v\ x))\ x$ **using** $add(5)\ integrable\text{-}simple\text{-}function[OF\ add(1,2)]\ integrable\text{-}simple\text{-}function[OF\ add(3,4)]$ **by** (*intro cond-exp-cong, auto*)

ultimately **show** *?case* **by** *force*

qed

lemma *has-cond-exp-simple-lim*:

fixes $f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology,\ banach\}$

assumes $integrable[measurable]:\ integrable\ M\ f$

and $\bigwedge i.\ simple\text{-}function\ M\ (s\ i)$

and $\bigwedge i.\ emeasure\ M\ \{y \in space\ M.\ s\ i\ y \neq 0\} \neq \infty$

and $\bigwedge x.\ x \in space\ M \implies (\lambda i.\ s\ i\ x) \longrightarrow f\ x$

and $\bigwedge x\ i.\ x \in space\ M \implies norm\ (s\ i\ x) \leq 2 * norm\ (f\ x)$

obtains r

where $has\text{-}cond\text{-}exp\ M\ F\ f\ (\lambda x.\ lim\ (\lambda i.\ cond\text{-}exp\ M\ F\ (s\ (r\ i))\ x))$

$AE\ x\ in\ M.\ convergent\ (\lambda i.\ cond\text{-}exp\ M\ F\ (s\ (r\ i))\ x)$

strict-mono r

proof –

have $[measurable]: (s\ i) \in \text{borel-measurable } M$ **for** i **using** $assms(2)$ **by** $(simp\ add: \text{borel-measurable-simple-function})$

have $integrable-s: integrable\ M\ (\lambda x. s\ i\ x)$ **for** i **using** $assms(2)$ $assms(3)$ $integrable-simple-function$ **by** $blast$

have $integrable-4f: integrable\ M\ (\lambda x. 4 * norm\ (f\ x))$ **using** $assms(1)$ **by** $simp$

have $integrable-2f: integrable\ M\ (\lambda x. 2 * norm\ (f\ x))$ **using** $assms(1)$ **by** $simp$

have $integrable-2-cond-exp-norm-f: integrable\ M\ (\lambda x. 2 * cond-exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x)$ **by** $fast$

have $emeasure\ M\ \{y \in space\ M. s\ i\ y - s\ j\ y \neq 0\} \leq emeasure\ M\ \{y \in space\ M. s\ i\ y \neq 0\} + emeasure\ M\ \{y \in space\ M. s\ j\ y \neq 0\}$ **for** $i\ j$ **using** $simple-functionD(2)[OF\ assms(2)]$ **by** $(intro\ order-trans[OF\ emeasure-mono\ emeasure-subadditive],\ auto)$

hence $fin-sup: emeasure\ M\ \{y \in space\ M. s\ i\ y - s\ j\ y \neq 0\} \neq \infty$ **for** $i\ j$ **using** $assms(3)$ **by** $(metis\ (mono-tags)\ ennreal-add-eq-top\ linorder-not-less\ top.not-eq-extremum\ infinity-ennreal-def)$

have $emeasure\ M\ \{y \in space\ M. norm\ (s\ i\ y - s\ j\ y) \neq 0\} \leq emeasure\ M\ \{y \in space\ M. s\ i\ y \neq 0\} + emeasure\ M\ \{y \in space\ M. s\ j\ y \neq 0\}$ **for** $i\ j$ **using** $simple-functionD(2)[OF\ assms(2)]$ **by** $(intro\ order-trans[OF\ emeasure-mono\ emeasure-subadditive],\ auto)$

hence $fin-sup-norm: emeasure\ M\ \{y \in space\ M. norm\ (s\ i\ y - s\ j\ y) \neq 0\} \neq \infty$ **for** $i\ j$ **using** $assms(3)$ **by** $(metis\ (mono-tags)\ ennreal-add-eq-top\ linorder-not-less\ top.not-eq-extremum\ infinity-ennreal-def)$

have $Cauchy: Cauchy\ (\lambda n. s\ n\ x)$ **if** $x \in space\ M$ **for** x **using** $assms(4)$ $LIM-SEQ-imp-Cauchy\ that$ **by** $blast$

hence $bounded-range-s: bounded\ (range\ (\lambda n. s\ n\ x))$ **if** $x \in space\ M$ **for** x **using** $that\ cauchy-imp-bounded$ **by** $fast$

have $AE\ x\ in\ M. (\lambda n. diameter\ \{s\ i\ x \mid i. n \leq i\}) \longrightarrow 0$ **using** $Cauchy\ cauchy-iff-diameter-tends-to-zero-and-bounded$ **by** $fast$

moreover **have** $(\lambda x. diameter\ \{s\ i\ x \mid i. n \leq i\}) \in \text{borel-measurable } M$ **for** n **using** $bounded-range-s\ \text{borel-measurable-diameter}$ **by** $measurable$

moreover **have** $AE\ x\ in\ M. norm\ (diameter\ \{s\ i\ x \mid i. n \leq i\}) \leq 4 * norm\ (f\ x)$ **for** n

proof –

{

fix x **assume** $x: x \in space\ M$

have $diameter\ \{s\ i\ x \mid i. n \leq i\} \leq 2 * norm\ (f\ x) + 2 * norm\ (f\ x)$ **by** $(intro\ diameter-le,\ blast,\ subst\ dist-norm[symmetric],\ intro\ dist-triangle3[THEN\ order-trans,\ of\ 0],\ intro\ add-mono)\ (auto\ intro: assms(5)[OF\ x])$

hence $norm\ (diameter\ \{s\ i\ x \mid i. n \leq i\}) \leq 4 * norm\ (f\ x)$ **using** $diameter-ge-0[OF\ bounded-subset[OF\ bounded-range-s],\ OF\ x,\ of\ \{s\ i\ x \mid i. n \leq i\}]$ **by** $force$

}

thus $?thesis$ **by** $fast$

qed

ultimately have *diameter-tendsto-zero*: $(\lambda n. \text{LINT } x|M. \text{diameter } \{s \ i \ x \mid i. n \leq i\}) \longrightarrow 0$ **by** (*intro integral-dominated-convergence*[*OF borel-measurable-const*[*of 0*] - *integrable-4f*, *simplified*]) (*fast+*)

have *diameter-integrable*: *integrable* $M (\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\})$ **for** n **using** *assms*(1,5) **by** (*intro integrable-bound-diameter*[*OF bounded-range-s integrable-2f*], *auto*)

have *dist-integrable*: *integrable* $M (\lambda x. \text{dist } (s \ i \ x) (s \ j \ x))$ **for** $i \ j$ **using** *assms*(5) *dist-triangle3*[*of s i - - 0*, *THEN order-trans*, *OF add-mono*, *of - 2 * norm (f -)*] **by** (*intro Bochner-Integration.integrable-bound*[*OF integrable-4f*]) *fastforce+*

hence *dist-norm-integrable*: *integrable* $M (\lambda x. \text{norm } (s \ i \ x - s \ j \ x))$ **for** $i \ j$ **unfolding** *dist-norm* **by** *presburger*

have $\exists N. \forall i \geq N. \forall j \geq N. \text{LINT } x|M. \text{norm } (\text{cond-exp } M \ F \ (s \ i) \ x - \text{cond-exp } M \ F \ (s \ j) \ x) < e$ **if** *e-pos*: $e > 0$ **for** e

proof -

obtain N **where** $*$: $\text{LINT } x|M. \text{diameter } \{s \ i \ x \mid i. n \leq i\} < e$ **if** $n \geq N$ **for** n **using** *that order-tendsto-iff*[*THEN iffD1*, *OF diameter-tendsto-zero*, *unfolded eventually-sequentially*] *e-pos* **by** *presburger*

{

fix $i \ j \ x$ **assume** *asm*: $i \geq N \ j \geq N \ x \in \text{space } M$

have *case-prod dist* ' $(\{s \ i \ x \mid i. N \leq i\} \times \{s \ i \ x \mid i. N \leq i\}) = \text{case-prod } (\lambda i \ j. \text{dist } (s \ i \ x) (s \ j \ x))$ ' $(\{N..\} \times \{N..\})$ **by** *fast*

hence *diameter* $\{s \ i \ x \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s \ i \ x) (s \ j \ x))$ **unfolding** *diameter-def* **by** *auto*

moreover have $(\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s \ i \ x) (s \ j \ x)) \geq \text{dist } (s \ i \ x) (s \ j \ x)$ **using** *asm bounded-imp-bdd-above*[*OF bounded-imp-dist-bounded*, *OF bounded-range-s*] **by** (*intro cSup-upper*, *auto*)

ultimately have *diameter* $\{s \ i \ x \mid i. N \leq i\} \geq \text{dist } (s \ i \ x) (s \ j \ x)$ **by** *presburger*

}

hence $\text{LINT } x|M. \text{dist } (s \ i \ x) (s \ j \ x) < e$ **if** $i \geq N \ j \geq N$ **for** $i \ j$ **using** *that ** **by** (*intro integral-mono*[*OF dist-integrable diameter-integrable*, *THEN order.strict-trans1*], *blast+*)

moreover have $\text{LINT } x|M. \text{norm } (\text{cond-exp } M \ F \ (s \ i) \ x - \text{cond-exp } M \ F \ (s \ j) \ x) \leq \text{LINT } x|M. \text{dist } (s \ i \ x) (s \ j \ x)$ **for** $i \ j$

proof -

have $\text{LINT } x|M. \text{norm } (\text{cond-exp } M \ F \ (s \ i) \ x - \text{cond-exp } M \ F \ (s \ j) \ x) = \text{LINT } x|M. \text{norm } (\text{cond-exp } M \ F \ (s \ i) \ x + -1 *_R \text{cond-exp } M \ F \ (s \ j) \ x)$ **unfolding** *dist-norm* **by** *simp*

also have $\dots = \text{LINT } x|M. \text{norm } (\text{cond-exp } M \ F \ (\lambda x. s \ i \ x - s \ j \ x) \ x)$ **using** *has-cond-exp-charact*(2)[*OF has-cond-exp-add*[*OF - has-cond-exp-scaleR-right*, *OF has-cond-exp-charact*(1,1), *OF has-cond-exp-simple*(1,1)[*OF assms*(2,3)]]], *THEN AE-symmetric*, *of i - 1 j*] **by** (*intro integral-cong-AE*) *force+*

also have $\dots \leq \text{LINT } x|M. \text{cond-exp } M \ F \ (\lambda x. \text{norm } (s \ i \ x - s \ j \ x)) \ x$ **using** *cond-exp-contraction-simple*[*OF - fin-sup*, *of i j*] *integrable-cond-exp assms*(2) **by**

(intro integral-mono-AE, fast+)

also have ... = $LINT\ x|M. norm\ (s\ i\ x - s\ j\ x)$ **unfolding** *set-integral-space*(1)[*OF integrable-cond-exp, symmetric*] *set-integral-space*[*OF dist-norm-integrable, symmetric*] **by** (intro *has-cond-expD*(1)[*OF has-cond-exp-simple*[*OF - fin-sup-norm*], *symmetric*]) (*metis assms*(2) *simple-function-compose1 simple-function-diff, metis sets.top subalg subalgebra-def*)

finally show ?thesis **unfolding** *dist-norm* .

qed

ultimately show ?thesis **using** *order.strict-trans1* **by** *meson*

qed

then obtain *r* **where** *strict-mono-r: strict-mono r* **and** *AE-Cauchy: AE x in M. Cauchy ($\lambda i. cond-exp\ M\ F\ (s\ (r\ i))\ x$)* **by** (*rule cauchy-L1-AE-cauchy-subseq*[*OF integrable-cond-exp*], *auto*)

hence *ae-lim-cond-exp: AE x in M. ($\lambda n. cond-exp\ M\ F\ (s\ (r\ n))\ x$) \longrightarrow lim ($\lambda n. cond-exp\ M\ F\ (s\ (r\ n))\ x$)* **using** *Cauchy-convergent-iff convergent-LIMSEQ-iff* **by** *fastforce*

have *cond-exp-bounded: AE x in M. norm (cond-exp M F (s (r n)) x) \leq cond-exp M F ($\lambda x. 2 * norm\ (f\ x)$) x* **for** *n*

proof –

have *AE x in M. norm (cond-exp M F (s (r n)) x) \leq cond-exp M F ($\lambda x. norm\ (s\ (r\ n)\ x)$) x* **by** (*rule cond-exp-contraction-simple*[*OF assms*(2,3)])

moreover have *AE x in M. real-cond-exp M F ($\lambda x. norm\ (s\ (r\ n)\ x)$) x \leq real-cond-exp M F ($\lambda x. 2 * norm\ (f\ x)$) x* **using** *integrable-s integrable-2f assms*(5) **by** (*intro real-cond-exp-mono, auto*)

ultimately show ?thesis **using** *cond-exp-real*[*OF integrable-norm, OF integrable-s, of r n*] *cond-exp-real*[*OF integrable-2f*] **by** *force*

qed

have *lim-integrable: integrable M ($\lambda x. lim\ (\lambda i. cond-exp\ M\ F\ (s\ (r\ i))\ x)$)* **by** (*intro integrable-dominated-convergence*[*OF - borel-measurable-cond-exp' integrable-cond-exp ae-lim-cond-exp cond-exp-bounded*], *simp*)

{

fix *A* **assume** *A-in-sets-F: A \in sets F*

have *AE x in M. norm (indicator A x *_R cond-exp M F (s (r n)) x) \leq cond-exp M F ($\lambda x. 2 * norm\ (f\ x)$) x* **for** *n*

proof –

have *AE x in M. norm (indicator A x *_R cond-exp M F (s (r n)) x) \leq norm (cond-exp M F (s (r n)) x)* **unfolding** *indicator-def* **by** *simp*

thus ?thesis **using** *cond-exp-bounded*[*of n*] **by** *force*

qed

hence *lim-cond-exp-int: ($\lambda n. LINT\ x:A|M. cond-exp\ M\ F\ (s\ (r\ n))\ x$) \longrightarrow LINT x:A|M. lim ($\lambda n. cond-exp\ M\ F\ (s\ (r\ n))\ x$)*

using *ae-lim-cond-exp measurable-from-subalg*[*OF subalg borel-measurable-indicator, OF A-in-sets-F*] *cond-exp-bounded*

unfolding *set-lebesgue-integral-def*

by (*intro integral-dominated-convergence*[*OF borel-measurable-scaleR borel-measurable-scaleR integrable-cond-exp*]) (*fastforce simp add: tendsto-scaleR*)+

have $AE\ x\ in\ M.\ norm\ (indicator\ A\ x\ *_R\ s\ (r\ n)\ x) \leq 2 * norm\ (f\ x)$ **for** n
proof –
have $AE\ x\ in\ M.\ norm\ (indicator\ A\ x\ *_R\ s\ (r\ n)\ x) \leq norm\ (s\ (r\ n)\ x)$
unfolding *indicator-def* **by** *simp*
thus *?thesis* **using** *assms(5)[of - r n]* **by** *fastforce*
qed
hence $lim-s-int: (\lambda n.\ LINT\ x:A|M.\ s\ (r\ n)\ x) \longrightarrow LINT\ x:A|M.\ f\ x$
using *measurable-from-subalg[OF subalg borel-measurable-indicator, OF A-in-sets-F]*
LIMSEQ-subseq-LIMSEQ[OF assms(4) strict-mono-r] *assms(5)*
unfolding *set-lebesgue-integral-def comp-def*
by (*intro integral-dominated-convergence[OF borel-measurable-scaleR borel-measurable-scaleR*
integrable-2f]) (*fastforce simp add: tendsto-scaleR*)**+**

have $LINT\ x:A|M.\ lim\ (\lambda n.\ cond-exp\ M\ F\ (s\ (r\ n))\ x) = lim\ (\lambda n.\ LINT\ x:A|M.\ cond-exp\ M\ F\ (s\ (r\ n))\ x)$ **using** *limI[OF lim-cond-exp-int]* **by** *argo*
also have $\dots = lim\ (\lambda n.\ LINT\ x:A|M.\ s\ (r\ n)\ x)$ **using** *has-cond-expD(1)[OF*
has-cond-exp-simple[OF assms(2,3)] A-in-sets-F, symmetric] **by** *presburger*
also have $\dots = LINT\ x:A|M.\ f\ x$ **using** *limI[OF lim-s-int]* **by** *argo*
finally have $LINT\ x:A|M.\ lim\ (\lambda n.\ cond-exp\ M\ F\ (s\ (r\ n))\ x) = LINT\ x:A|M.\ f\ x$.
}
hence $has-cond-exp\ M\ F\ f\ (\lambda x.\ lim\ (\lambda i.\ cond-exp\ M\ F\ (s\ (r\ i))\ x))$ **using**
assms(1) lim-integrable **by** (*intro has-cond-expI', auto*)
thus *thesis* **using** *AE-Cauchy Cauchy-convergent strict-mono-r* **by** (*auto intro!: that*)
qed

lemma *cond-exp-simple-lim*:

fixes $f :: 'a \Rightarrow 'b :: \{second-countable-topology, banach\}$
assumes *[measurable]:integrable M f*
and $\bigwedge i.\ simple-function\ M\ (s\ i)$
and $\bigwedge i.\ emeasure\ M\ \{y \in space\ M.\ s\ i\ y \neq 0\} \neq \infty$
and $\bigwedge x.\ x \in space\ M \implies (\lambda i.\ s\ i\ x) \longrightarrow f\ x$
and $\bigwedge x\ i.\ x \in space\ M \implies norm\ (s\ i\ x) \leq 2 * norm\ (f\ x)$
obtains $r\ where\ AE\ x\ in\ M.\ (\lambda i.\ cond-exp\ M\ F\ (s\ (r\ i))\ x) \longrightarrow cond-exp\ M\ F\ f\ x\ strict-mono\ r$
proof –
obtain $r\ where\ AE\ x\ in\ M.\ cond-exp\ M\ F\ f\ x = lim\ (\lambda i.\ cond-exp\ M\ F\ (s\ (r\ i))\ x)$ $AE\ x\ in\ M.\ convergent\ (\lambda i.\ cond-exp\ M\ F\ (s\ (r\ i))\ x)\ strict-mono\ r$ **using**
has-cond-exp-charact(2) **by** (*auto intro: has-cond-exp-simple-lim[OF assms]*)
thus *?thesis* **by** (*auto intro!: that[of r] simp: convergent-LIMSEQ-iff*)
qed

corollary *has-cond-expI*:

fixes $f :: 'a \Rightarrow 'b :: \{second-countable-topology, banach\}$
assumes *integrable M f*
shows $has-cond-exp\ M\ F\ f\ (cond-exp\ M\ F\ f)$
proof –
obtain $s\ where\ s-is: \bigwedge i.\ simple-function\ M\ (s\ i) \bigwedge i.\ emeasure\ M\ \{y \in space\ M.$

$s \ i \ y \neq 0\} \neq \infty \wedge x. x \in \text{space } M \implies (\lambda i. s \ i \ x) \longrightarrow f \ x \wedge x \ i. x \in \text{space } M \implies$
 $\text{norm } (s \ i \ x) \leq 2 * \text{norm } (f \ x)$ **using** *integrable-implies-simple-function-sequence*[*OF*
assms] **by** *blast*
show *?thesis* **using** *has-cond-exp-simple-lim*[*OF* *assms* *s-is*] *has-cond-exp-charact*(1)
by *metis*
qed

lemma *cond-exp-nested-subalg*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M \ f$ *subalgebra* $M \ G$ *subalgebra* $G \ F$
shows $AE \ \xi \text{ in } M. \text{cond-exp } M \ F \ f \ \xi = \text{cond-exp } M \ F \ (\text{cond-exp } M \ G \ f) \ \xi$
using *has-cond-expI* *assms* *sigma-finite-subalgebra-def* **by** (*auto intro!*: *has-cond-exp-nested-subalg*[*THEN*
has-cond-exp-charact(2), *THEN* *AE-symmetric*] *sigma-finite-subalgebra.has-cond-expI*[*OF*
sigma-finite-subalgebra.intro[*OF* *assms*(2)]] *nested-subalg-is-sigma-finite*)

lemma *cond-exp-set-integral*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M \ f \ A \in \text{sets } F$
shows $(\int x \in A. f \ x \ \partial M) = (\int x \in A. \text{cond-exp } M \ F \ f \ x \ \partial M)$
using *has-cond-expD*(1)[*OF* *has-cond-expI*, *OF* *assms*] **by** *argo*

lemma *cond-exp-add*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M \ f$ *integrable* $M \ g$
shows $AE \ x \text{ in } M. \text{cond-exp } M \ F \ (\lambda x. f \ x + g \ x) \ x = \text{cond-exp } M \ F \ f \ x +$
 $\text{cond-exp } M \ F \ g \ x$
using *has-cond-exp-add*[*OF* *has-cond-expI*(1,1), *OF* *assms*, *THEN* *has-cond-exp-charact*(2)]
.

lemma *cond-exp-diff*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M \ f$ *integrable* $M \ g$
shows $AE \ x \text{ in } M. \text{cond-exp } M \ F \ (\lambda x. f \ x - g \ x) \ x = \text{cond-exp } M \ F \ f \ x -$
 $\text{cond-exp } M \ F \ g \ x$
using *has-cond-exp-add*[*OF* - *has-cond-exp-scaleR-right*, *OF* *has-cond-expI*(1,1),
OF *assms*, *THEN* *has-cond-exp-charact*(2), *of* -1] **by** *simp*

lemma *cond-exp-diff'*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M \ f$ *integrable* $M \ g$
shows $AE \ x \text{ in } M. \text{cond-exp } M \ F \ (f - g) \ x = \text{cond-exp } M \ F \ f \ x - \text{cond-exp } M$
 $F \ g \ x$
unfolding *fun-diff-def* **using** *assms* **by** (*rule* *cond-exp-diff*)

lemma *cond-exp-scaleR-left*:

fixes $f :: 'a \Rightarrow \text{real}$
assumes *integrable* $M \ f$

shows $AE\ x\ in\ M. \text{cond-exp } M\ F\ (\lambda x. f\ x *_{\mathbb{R}} c)\ x = \text{cond-exp } M\ F\ f\ x *_{\mathbb{R}} c$
using $\text{cond-exp-set-integral}[OF\ \text{assms}]\ \text{subalg}\ \text{assms}$ **unfolding** subalgebra-def
by $(\text{intro}\ \text{cond-exp-charact},$
 $\text{subst}\ \text{set-integral-scaleR-left}, \text{blast}, \text{intro}\ \text{assms},$
 $\text{subst}\ \text{set-integral-scaleR-left}, \text{blast}, \text{intro}\ \text{integrable-cond-exp})$
 auto

lemma $\text{cond-exp-contraction}$:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$

assumes $\text{integrable } M\ f$

shows $AE\ x\ in\ M. \text{norm } (\text{cond-exp } M\ F\ f\ x) \leq \text{cond-exp } M\ F\ (\lambda x. \text{norm } (f\ x))$

x

proof –

obtain s **where** $s: \bigwedge i. \text{simple-function } M\ (s\ i) \bigwedge i. \text{emeasure } M\ \{y \in \text{space } M. s\ i\ y \neq 0\} \neq \infty \bigwedge x. x \in \text{space } M \implies (\lambda i. s\ i\ x) \longrightarrow f\ x \bigwedge i. x \in \text{space } M \implies \text{norm } (s\ i\ x) \leq 2 * \text{norm } (f\ x)$

by $(\text{blast}\ \text{intro: integrable-implies-simple-function-sequence}[OF\ \text{assms}])$

obtain r **where** $r: AE\ x\ in\ M. (\lambda i. \text{cond-exp } M\ F\ (s\ (r\ i))\ x) \longrightarrow \text{cond-exp } M\ F\ f\ x \text{ strict-mono } r$ **using** $\text{cond-exp-simple-lim}[OF\ \text{assms } s]$ **by** blast

have $\text{norm-s-r}: \bigwedge i. \text{simple-function } M\ (\lambda x. \text{norm } (s\ (r\ i)\ x)) \bigwedge i. \text{emeasure } M\ \{y \in \text{space } M. \text{norm } (s\ (r\ i)\ y) \neq 0\} \neq \infty \bigwedge x. x \in \text{space } M \implies (\lambda i. \text{norm } (s\ (r\ i)\ x)) \longrightarrow \text{norm } (f\ x) \bigwedge i. x \in \text{space } M \implies \text{norm } (\text{norm } (s\ (r\ i)\ x)) \leq 2 * \text{norm } (\text{norm } (f\ x))$

using s **by** $(\text{auto}\ \text{intro: LIMSEQ-subseq-LIMSEQ}[OF\ \text{tendsto-norm } r(2), \text{unfolded comp-def}]\ \text{simple-function-compose1})$

obtain r' **where** $r': AE\ x\ in\ M. (\lambda i. (\text{cond-exp } M\ F\ (\lambda x. \text{norm } (s\ (r\ (r'\ i))\ x))\ x)) \longrightarrow \text{cond-exp } M\ F\ (\lambda x. \text{norm } (f\ x))\ x \text{ strict-mono } r'$ **using** $\text{cond-exp-simple-lim}[OF\ \text{integrable-norm norm-s-r}, OF\ \text{assms}]$ **by** blast

have $AE\ x\ in\ M. \forall i. \text{norm } (\text{cond-exp } M\ F\ (s\ (r\ (r'\ i))\ x)) \leq \text{cond-exp } M\ F\ (\lambda x. \text{norm } (s\ (r\ (r'\ i))\ x))\ x$ **using** s **by** $(\text{auto}\ \text{intro: cond-exp-contraction-simple simp add: AE-all-countable})$

moreover **have** $AE\ x\ in\ M. (\lambda i. \text{norm } (\text{cond-exp } M\ F\ (s\ (r\ (r'\ i))\ x)) \longrightarrow \text{norm } (\text{cond-exp } M\ F\ f\ x))$ **using** r $\text{LIMSEQ-subseq-LIMSEQ}[OF\ \text{tendsto-norm } r'(2), \text{unfolded comp-def}]$ **by** fast

ultimately show $?thesis$ **using** $\text{LIMSEQ-le } r'(1)$ **by** fast

qed

lemma $\text{cond-exp-measurable-mult}$:

fixes $f\ g :: 'a \Rightarrow \text{real}$

assumes $[\text{measurable}]: \text{integrable } M\ (\lambda x. f\ x * g\ x) \text{ integrable } M\ g\ f \in \text{borel-measurable } F$

shows $\text{integrable } M\ (\lambda x. f\ x * \text{cond-exp } M\ F\ g\ x)$

$AE\ x\ in\ M. \text{cond-exp } M\ F\ (\lambda x. f\ x * g\ x)\ x = f\ x * \text{cond-exp } M\ F\ g\ x$

proof –

show *integrable*: *integrable* M $(\lambda x. f x * \text{cond-exp } M F g x)$ **using** *cond-exp-real*[*OF* *assms*(2)] **by** (*intro integrable-cong-AE-imp*[*OF* *real-cond-exp-intg*(1), *OF* *assms*(1,3) *assms*(2)[*THEN borel-measurable-integrable*]] *measurable-from-subalg*[*OF* *subalg*]) *auto*

interpret *sigma-finite-measure restr-to-subalg* $M F$ **by** (*rule sigma-fin-subalg*)

{

fix A **assume** *asm*: $A \in \text{sets } F$

hence *asm'*: $A \in \text{sets } M$ **using** *subalg* **by** (*fastforce simp add: subalgebra-def*)

have *set-lebesgue-integral* $M A$ (*cond-exp* $M F$ $(\lambda x. f x * g x)$) = *set-lebesgue-integral* $M A$ $(\lambda x. f x * g x)$ **by** (*simp add: cond-exp-set-integral*[*OF* *assms*(1) *asm*])

also have ... = *set-lebesgue-integral* $M A$ $(\lambda x. f x * \text{real-cond-exp } M F g x)$ **using** *borel-measurable-times*[*OF* *borel-measurable-indicator*[*OF* *asm*] *assms*(3)] *borel-measurable-integrable*[*OF* *assms*(2)] *integrable-mult-indicator*[*OF* *asm'* *assms*(1)] **by** (*fastforce simp add: set-lebesgue-integral-def mult.assoc[symmetric] intro: real-cond-exp-intg(2)[symmetric]*)

also have ... = *set-lebesgue-integral* $M A$ $(\lambda x. f x * \text{cond-exp } M F g x)$ **using** *cond-exp-real*[*OF* *assms*(2)] *asm'* *borel-measurable-cond-exp'* *borel-measurable-cond-exp2* *measurable-from-subalg*[*OF* *subalg* *assms*(3)] **by** (*auto simp add: set-lebesgue-integral-def intro: integral-cong-AE*)

finally have *set-lebesgue-integral* $M A$ (*cond-exp* $M F$ $(\lambda x. f x * g x)$) = $\int_{x \in A. (f x * \text{cond-exp } M F g x) \partial M}$.

}

hence $AE x \text{ in } \text{restr-to-subalg } M F. \text{cond-exp } M F (\lambda x. f x * g x) x = f x * \text{cond-exp } M F g x$ **by** (*intro density-unique integrable-cond-exp integrable integrable-in-subalg subalg, measurable, simp add: set-lebesgue-integral-def integral-subalgebra2*[*OF* *subalg*] *sets-restr-to-subalg*[*OF* *subalg*])

thus $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x * g x) x = f x * \text{cond-exp } M F g x$ **by** (*rule AE-restr-to-subalg*[*OF* *subalg*])

qed

lemma *cond-exp-measurable-scaleR*:

fixes $f :: 'a \Rightarrow \text{real}$ **and** $g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$

assumes [*measurable*]: *integrable* M $(\lambda x. f x *_R g x)$ *integrable* $M g f \in \text{borel-measurable } F$

shows *integrable* M $(\lambda x. f x *_R \text{cond-exp } M F g x)$

$AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x *_R g x) x = f x *_R \text{cond-exp } M F g x$

proof –

let $?F = \text{restr-to-subalg } M F$

have *subalg'*: *subalgebra* M (*restr-to-subalg* $M F$) **by** (*metis sets-eq-imp-space-eq sets-restr-to-subalg subalg subalgebra-def*)

{

fix z **assume** *asm*[*measurable*]: *integrable* M $(\lambda x. z x *_R g x)$ $z \in \text{borel-measurable } ?F$

hence *asm'*[*measurable*]: $z \in \text{borel-measurable } F$ **using** *measurable-in-subalg'* *subalg* **by** *blast*

have *integrable* M $(\lambda x. z x *_R \text{cond-exp } M F g x)$ *LINT* $x | M. z x *_R g x = \text{LINT } x | M. z x *_R \text{cond-exp } M F g x$

proof –

obtain s **where** s -is: $\bigwedge i. \text{simple-function } ?F (s i) \bigwedge x. x \in \text{space } ?F \implies (\lambda i.$

$s \ i \ x) \longrightarrow z \ x \ \bigwedge i \ x. x \in \text{space } ?F \implies \text{norm } (s \ i \ x) \leq 2 * \text{norm } (z \ x)$ **using** *borel-measurable-implies-sequence-metric*[*OF asm(2), of 0*] **by force**

have *s-scaleR-g-tendsto*: $AE \ x \ \text{in } M. (\lambda i. s \ i \ x *_{\mathbb{R}} g \ x) \longrightarrow z \ x *_{\mathbb{R}} g \ x$ **using** *s-is(2)* **by** (*simp add: space-restr-to-subalg tendsto-scaleR*)

have *s-scaleR-cond-exp-g-tendsto*: $AE \ x \ \text{in } ?F. (\lambda i. s \ i \ x *_{\mathbb{R}} \text{cond-exp } M \ F \ g \ x) \longrightarrow z \ x *_{\mathbb{R}} \text{cond-exp } M \ F \ g \ x$ **using** *s-is(2)* **by** (*simp add: tendsto-scaleR*)

have *s-scaleR-g-meas*: $(\lambda x. s \ i \ x *_{\mathbb{R}} g \ x) \in \text{borel-measurable } M$ **for** *i* **using** *s-is(1)*[*THEN borel-measurable-simple-function, THEN subalg'*[*THEN measurable-from-subalg*]] **by simp**

have *s-scaleR-cond-exp-g-meas*: $(\lambda x. s \ i \ x *_{\mathbb{R}} \text{cond-exp } M \ F \ g \ x) \in \text{borel-measurable } ?F$ **for** *i* **using** *s-is(1)*[*THEN borel-measurable-simple-function*] *measurable-in-subalg*[*OF subalg borel-measurable-cond-exp*] **by** (*fastforce intro: borel-measurable-scaleR*)

have *s-scaleR-g-AE-bdd*: $AE \ x \ \text{in } M. \text{norm } (s \ i \ x *_{\mathbb{R}} g \ x) \leq 2 * \text{norm } (z \ x *_{\mathbb{R}} g \ x)$ **for** *i* **using** *s-is(3)* **by** (*fastforce simp add: space-restr-to-subalg mult.assoc[symmetric] mult-right-mono*)

{

fix *i*

have *asm: integrable M* $(\lambda x. \text{norm } (z \ x) * \text{norm } (g \ x))$ **using** *asm(1)*[*THEN integrable-norm*] **by simp**

have $AE \ x \ \text{in } ?F. \text{norm } (s \ i \ x *_{\mathbb{R}} \text{cond-exp } M \ F \ g \ x) \leq 2 * \text{norm } (z \ x) * \text{norm } (\text{cond-exp } M \ F \ g \ x)$ **using** *s-is(3)* **by** (*fastforce simp add: mult-mono*)

moreover have $AE \ x \ \text{in } ?F. \text{norm } (z \ x) * \text{cond-exp } M \ F \ (\lambda x. \text{norm } (g \ x)) \ x = \text{cond-exp } M \ F \ (\lambda x. \text{norm } (z \ x) * \text{norm } (g \ x)) \ x$ **by** (*rule cond-exp-measurable-mult(2)*[*THEN AE-symmetric, OF asm integrable-norm, OF assms(2), THEN AE-restr-to-subalg2*[*OF subalg*]], *auto*)

ultimately have $AE \ x \ \text{in } ?F. \text{norm } (s \ i \ x *_{\mathbb{R}} \text{cond-exp } M \ F \ g \ x) \leq 2 * \text{cond-exp } M \ F \ (\lambda x. \text{norm } (z \ x *_{\mathbb{R}} g \ x)) \ x$ **using** *cond-exp-contraction*[*OF assms(2), THEN AE-restr-to-subalg2*[*OF subalg*]] *order-trans*[*OF - mult-mono*] **by fastforce**

}

note *s-scaleR-cond-exp-g-AE-bdd = this*

{

fix *i*

have *s-meas-M[measurable]*: $s \ i \in \text{borel-measurable } M$ **by** (*meson borel-measurable-simple-function measurable-from-subalg s-is(1) subalg'*)

have *s-meas-F[measurable]*: $s \ i \in \text{borel-measurable } F$ **by** (*meson borel-measurable-simple-function measurable-in-subalg' s-is(1) subalg*)

have *s-scaleR-eq*: $s \ i \ x *_{\mathbb{R}} h \ x = (\sum_{y \in s \ i} \text{'space } M. (\text{indicator } (s \ i - \{y\} \cap \text{space } M) \ x *_{\mathbb{R}} y) *_{\mathbb{R}} h \ x)$ **if** $x \in \text{space } M$ **for** x **and** $h :: 'a \Rightarrow 'b$ **using** *simple-function-indicator-representation*[*OF s-is(1), of x i*] *that unfolding space-restr-to-subalg scaleR-left.sum*[*of - - h x, symmetric*] **by presburger**

have $LINT\ x|M. s\ i\ x *_R g\ x = LINT\ x|M. (\sum_{y \in s\ i\ ' \text{ space } M. indicator\ (s\ i\ -'\ \{y\} \cap \text{space } M)\ x *_R y *_R g\ x})$ **using** $s\text{-scaleR-eq}$ **by** (intro Bochner-Integration.integral-cong) auto
also have $\dots = (\sum_{y \in s\ i\ ' \text{ space } M. LINT\ x|M. indicator\ (s\ i\ -'\ \{y\} \cap \text{space } M)\ x *_R y *_R g\ x})$ **by** (intro Bochner-Integration.integral-sum integrable-mult-indicator[OF - integrable-scaleR-right] assms(2)) simp
also have $\dots = (\sum_{y \in s\ i\ ' \text{ space } M. y *_R set\text{-lebesgue-integral } M\ (s\ i\ -'\ \{y\} \cap \text{space } M)\ g})$ **by** (simp only: set-lebesgue-integral-def[symmetric]) simp
also have $\dots = (\sum_{y \in s\ i\ ' \text{ space } M. y *_R set\text{-lebesgue-integral } M\ (s\ i\ -'\ \{y\} \cap \text{space } M)\ (cond\text{-exp } M\ F\ g)})$ **using** assms(2) subalg borel-measurable-vimage[OF s-meas-F] **by** (subst cond-exp-set-integral, auto simp add: subalgebra-def)
also have $\dots = (\sum_{y \in s\ i\ ' \text{ space } M. LINT\ x|M. indicator\ (s\ i\ -'\ \{y\} \cap \text{space } M)\ x *_R y *_R cond\text{-exp } M\ F\ g\ x})$ **by** (simp only: set-lebesgue-integral-def[symmetric]) simp
also have $\dots = LINT\ x|M. (\sum_{y \in s\ i\ ' \text{ space } M. indicator\ (s\ i\ -'\ \{y\} \cap \text{space } M)\ x *_R y *_R cond\text{-exp } M\ F\ g\ x})$ **by** (intro Bochner-Integration.integral-sum[symmetric] integrable-mult-indicator[OF - integrable-scaleR-right]) auto
also have $\dots = LINT\ x|M. s\ i\ x *_R cond\text{-exp } M\ F\ g\ x$ **using** $s\text{-scaleR-eq}$ **by** (intro Bochner-Integration.integral-cong) auto
finally have $LINT\ x|M. s\ i\ x *_R g\ x = LINT\ x|?F. s\ i\ x *_R cond\text{-exp } M\ F\ g\ x$ **by** (simp add: integral-subalgebra2[OF subalg])
}
note $integral\text{-s-eq} = this$

show $integrable\ M\ (\lambda x. z\ x *_R cond\text{-exp } M\ F\ g\ x)$ **using** $s\text{-scaleR-cond-exp-g-meas}\ asm(2)\ borel\text{-measurable-cond-exp}'$ **by** (intro integrable-from-subalg[OF subalg] integrable-cond-exp integrable-dominated-convergence[OF - - - $s\text{-scaleR-cond-exp-g-tendsto}\ s\text{-scaleR-cond-exp-g-AE-bdd}$]) (auto intro: measurable-from-subalg[OF subalg] integrable-in-subalg measurable-in-subalg subalg)

have $(\lambda i. LINT\ x|M. s\ i\ x *_R g\ x) \longrightarrow LINT\ x|M. z\ x *_R g\ x$ **using** $s\text{-scaleR-g-meas}\ asm(1)[THEN\ integrable\text{-norm}]\ asm'\ borel\text{-measurable-cond-exp}'$ **by** (intro integral-dominated-convergence[OF - - - $s\text{-scaleR-g-tendsto}\ s\text{-scaleR-g-AE-bdd}$]) (auto intro: measurable-from-subalg[OF subalg])

moreover have $(\lambda i. LINT\ x|?F. s\ i\ x *_R cond\text{-exp } M\ F\ g\ x) \longrightarrow LINT\ x|?F. z\ x *_R cond\text{-exp } M\ F\ g\ x$ **using** $s\text{-scaleR-cond-exp-g-meas}\ asm(2)\ borel\text{-measurable-cond-exp}'$ **by** (intro integral-dominated-convergence[OF - - - $s\text{-scaleR-cond-exp-g-tendsto}\ s\text{-scaleR-cond-exp-g-AE-bdd}$]) (auto intro: measurable-from-subalg[OF subalg] integrable-in-subalg measurable-in-subalg subalg)

ultimately show $LINT\ x|M. z\ x *_R g\ x = LINT\ x|M. z\ x *_R cond\text{-exp } M\ F\ g\ x$ **using** $integral\text{-s-eq}$ **using** subalg **by** (simp add: LIMSEQ-unique integral-subalgebra2)

qed

}

note $* = this$

show *integrable* M $(\lambda x. f x *_R \text{cond-exp } M F g x)$ **using** $* \text{assms measurable-in-subalg}[OF \text{subalg}]$ **by** *blast*

```
{
  fix A assume asm: A ∈ F
  hence integrable M (λx. indicat-real A x *_R f x *_R g x) using subalg by
  (fastforce simp add: subalgebra-def intro!: integrable-mult-indicator assms(1))
  hence set-lebesgue-integral M A (λx. f x *_R g x) = set-lebesgue-integral M A
  (λx. f x *_R cond-exp M F g x) unfolding set-lebesgue-integral-def using asm by
  (auto intro!: * measurable-in-subalg[OF subalg])
}
thus AE x in M. cond-exp M F (λx. f x *_R g x) x = f x *_R cond-exp M F g x
using borel-measurable-cond-exp by (intro cond-exp-charact, auto intro!: * assms
measurable-in-subalg[OF subalg])
qed
```

lemma *cond-exp-sum* [intro, simp]:

```
fixes f :: 't ⇒ 'a ⇒ 'b :: {second-countable-topology, banach}
assumes [measurable]: ∧i. integrable M (f i)
shows AE x in M. cond-exp M F (λx. ∑ i∈I. f i x) x = (∑ i∈I. cond-exp M F
(f i) x)
proof (rule has-cond-exp-charact, intro has-cond-expI')
  fix A assume [measurable]: A ∈ sets F
  then have A-meas [measurable]: A ∈ sets M by (meson subsetD subalg subalge-
bra-def)
```

```
  have (∫ x∈A. (∑ i∈I. f i x) ∂M) = (∫ x. (∑ i∈I. indicator A x *_R f i x) ∂M)
  unfolding set-lebesgue-integral-def by (simp add: scaleR-sum-right)
  also have ... = (∫ x. indicator A x *_R (∑ i∈I. f i x) ∂M) using assms by (auto
  intro!: Bochner-Integration.integral-sum integrable-mult-indicator)
  also have ... = (∫ x. indicator A x *_R cond-exp M F (f i) x ∂M) using
  cond-exp-set-integral[OF assms] by (simp add: set-lebesgue-integral-def)
  also have ... = (∫ x. (∑ i∈I. indicator A x *_R cond-exp M F (f i) x) ∂M)
  using assms by (auto intro!: Bochner-Integration.integral-sum[symmetric] inte-
  grable-mult-indicator)
  also have ... = (∫ x∈A. (∑ i∈I. cond-exp M F (f i) x) ∂M) unfolding set-lebesgue-integral-def
  by (simp add: scaleR-sum-right)
  finally show (∫ x∈A. (∑ i∈I. f i x) ∂M) = (∫ x∈A. (∑ i∈I. cond-exp M F (f i)
  x) ∂M) by auto
qed (auto simp add: assms integrable-cond-exp)
```

0.7 Ordered Banach Spaces

lemma *cond-exp-gr-c*:

```
fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
dered-real-vector}
assumes integrable M f AE x in M. f x > c
```

shows $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ f\ x > c$
proof –
define X **where** $X = \{x \in space\ M.\ cond\text{-}exp\ M\ F\ f\ x \leq c\}$
have $[measurable]: X \in sets\ F$ **unfolding** $X\text{-}def$ **by** $measurable\ (metis\ sets.top\ subalg\ subalgebra\text{-}def)$
hence $X\text{-}in\text{-}M: X \in sets\ M$ **using** $sets\text{-}restr\text{-}to\text{-}subalg\ subalg\ subalgebra\text{-}def$ **by** $blast$
have $emeasure\ M\ X = 0$
proof $(rule\ ccontr)$
assume $emeasure\ M\ X \neq 0$
have $emeasure\ (restr\text{-}to\text{-}subalg\ M\ F)\ X = emeasure\ M\ X$ **by** $(simp\ add: emeasure\text{-}restr\text{-}to\text{-}subalg\ subalg)$
hence $emeasure\ (restr\text{-}to\text{-}subalg\ M\ F)\ X > 0$ **using** $\neg(emeasure\ M\ X) = 0$ **by** $gr\text{-}zeroI$ **by** $auto$
then obtain A **where** $A: A \in sets\ (restr\text{-}to\text{-}subalg\ M\ F)\ A \subseteq X$ $emeasure\ (restr\text{-}to\text{-}subalg\ M\ F)\ A > 0$ $emeasure\ (restr\text{-}to\text{-}subalg\ M\ F)\ A < \infty$
using $\sigma\text{-}fin\text{-}subalg$ **by** $(metis\ emeasure\text{-}notin\text{-}sets\ ennreal\text{-}0\ infinity\text{-}ennreal\text{-}def\ le\text{-}less\text{-}linear\ neq\text{-}top\text{-}trans\ not\text{-}gr\text{-}zero\ order\text{-}reft\ \sigma\text{-}finite\text{-}measure.\ approx\text{-}PInf\text{-}emeasure\text{-}with\text{-}finite)$
hence $[simp]: A \in sets\ F$ **using** $subalg\ sets\text{-}restr\text{-}to\text{-}subalg$ **by** $blast$
hence $A\text{-}in\text{-}sets\text{-}M[simp]: A \in sets\ M$ **using** $sets\text{-}restr\text{-}to\text{-}subalg\ subalg\ subalgebra\text{-}def$ **by** $blast$
have $[simp]: set\text{-}integrable\ M\ A\ (\lambda x.\ c)$ **using** $A\ subalg$ **by** $(auto\ simp\ add: set\text{-}integrable\text{-}def\ emeasure\text{-}restr\text{-}to\text{-}subalg)$
have $[simp]: set\text{-}integrable\ M\ A\ f$ **unfolding** $set\text{-}integrable\text{-}def$ **by** $(rule\ integrable\text{-}mult\text{-}indicator,\ auto\ simp\ add: assms(1))$
have $AE\ x\ in\ M.\ indicator\ A\ x\ *_R\ c = indicator\ A\ x\ *_R\ f\ x$
proof $(rule\ integral\text{-}eq\text{-}mono\text{-}AE\text{-}eq\text{-}AE)$
show $LINT\ x|M.\ indicator\ A\ x\ *_R\ c = LINT\ x|M.\ indicator\ A\ x\ *_R\ f\ x$
proof $(simp\ only: set\text{-}lebesgue\text{-}integral\text{-}def[symmetric], rule\ antisym)$
show $(\int x \in A.\ c\ \partial M) \leq (\int x \in A.\ f\ x\ \partial M)$ **using** $assms(2)$ **by** $(intro\ set\text{-}integral\text{-}mono\text{-}AE\text{-}banach)\ auto$
have $(\int x \in A.\ f\ x\ \partial M) = (\int x \in A.\ cond\text{-}exp\ M\ F\ f\ x\ \partial M)$ **by** $(rule\ cond\text{-}exp\text{-}set\text{-}integral,\ auto\ simp\ add: assms)$
also have $\dots \leq (\int x \in A.\ c\ \partial M)$ **using** A **by** $(auto\ intro!: set\text{-}integral\text{-}mono\text{-}banach\ simp\ add: X\text{-}def)$
finally show $(\int x \in A.\ f\ x\ \partial M) \leq (\int x \in A.\ c\ \partial M)$ **by** $simp$
qed
show $AE\ x\ in\ M.\ indicator\ A\ x\ *_R\ c \leq indicator\ A\ x\ *_R\ f\ x$ **using** $assms$ **by** $(auto\ simp\ add: X\text{-}def\ indicator\text{-}def)$
qed $(auto\ simp\ add: set\text{-}integrable\text{-}def[symmetric])$
hence $AE\ x \in A\ in\ M.\ c = f\ x$ **by** $auto$
hence $AE\ x \in A\ in\ M.\ False$ **using** $assms(2)$ **by** $auto$
hence $A \in null\text{-}sets\ M$ **using** $AE\text{-}iff\text{-}null\text{-}sets\ A\text{-}in\text{-}sets\text{-}M$ **by** $metis$
thus $False$ **using** $A(3)$ **by** $(simp\ add: emeasure\text{-}restr\text{-}to\text{-}subalg\ null\text{-}setsD1\ subalg)$
qed
thus $?thesis$ **using** $AE\text{-}iff\text{-}null\text{-}sets[OF\ X\text{-}in\text{-}M]$ **unfolding** $X\text{-}def$ **by** $auto$
qed

corollary *cond-exp-less-c*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes *integrable* $M f$ *AE* x *in* M . $f x < c$
shows *AE* x *in* M . *cond-exp* $M F f x < c$
proof –
have *AE* x *in* M . *cond-exp* $M F f x = - \text{cond-exp } M F (\lambda x. - f x) x$ **using** *cond-exp-uminus*[*OF* *assms*(1)] **by** *auto*
moreover **have** *AE* x *in* M . *cond-exp* $M F (\lambda x. - f x) x > - c$ **using** *assms*
by (*intro cond-exp-gr-c*) *auto*
ultimately show *?thesis* **by** (*force simp add: minus-less-iff*)
qed

lemma *cond-exp-mono-strict*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes *integrable* $M f$ *integrable* $M g$ *AE* x *in* M . $f x < g x$
shows *AE* x *in* M . *cond-exp* $M F f x < \text{cond-exp } M F g x$
using *cond-exp-less-c*[*OF* *Bochner-Integration.integrable-diff*, *OF* *assms*(1,2), *of* 0]
cond-exp-diff[*OF* *assms*(1,2)] *assms*(3) **by** *auto*

lemma *cond-exp-ge-c*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes [*measurable*]: *integrable* $M f$
and *AE* x *in* M . $f x \geq c$
shows *AE* x *in* M . *cond-exp* $M F f x \geq c$
proof –
let $?F = \text{restr-to-subalg } M F$
interpret *sigma-finite-measure* *restr-to-subalg* $M F$ **using** *sigma-fin-subalg* **by** *auto*
{
fix A **assume** *asm*: $A \in \text{sets } ?F$ $0 < \text{measure } ?F A$
have [*simp*]: $\text{sets } ?F = \text{sets } F$ $\text{measure } ?F A = \text{measure } M A$ **using** *asm* **by** (*auto simp add: measure-def sets-restr-to-subalg*[*OF* *subalg*] *emeasure-restr-to-subalg*[*OF* *subalg*])
have $M-A$: $\text{emeasure } M A < \infty$ **using** *measure-zero-top* *asm* **by** (*force simp add: top.not-eq-extremum*)
hence $F-A$: $\text{emeasure } ?F A < \infty$ **using** *asm*(1) *emeasure-restr-to-subalg* *subalg* **by** *fastforce*
have *set-lebesgue-integral* $M A (\lambda-. c) \leq \text{set-lebesgue-integral } M A f$ **using** *assms* *asm* $M-A$ *subalg* **by** (*intro set-integral-mono-AE-banach*, *auto simp add: set-integrable-def integrable-mult-indicator subalgebra-def sets-restr-to-subalg*)
also **have** $\dots = \text{set-lebesgue-integral } M A (\text{cond-exp } M F f)$ **using** *cond-exp-set-integral*[*OF* *assms*(1)] *asm* **by** *auto*
also **have** $\dots = \text{set-lebesgue-integral } ?F A (\text{cond-exp } M F f)$ **unfolding** *set-lebesgue-integral-def* **using** *asm* *borel-measurable-cond-exp* **by** (*intro integral-subalgebra2*[*OF* *subalg*, *sym-metric*], *simp*)

finally have $(1 / \text{measure } ?F A) *_{\mathbb{R}} \text{set-lebesgue-integral } ?F A (\text{cond-exp } M F f)$
 $\in \{c..\}$ **using** *asm subalg M-A* **by** (*auto simp add: set-integral-const subalgebra-def*
intro!: pos-divideR-le-eq[THEN iffD1])
}
thus *?thesis* **using** *AE-restr-to-subalg[OF subalg] averaging-theorem[OF inte-*
grable-in-subalg closed-atLeast, OF subalg borel-measurable-cond-exp integrable-cond-exp]
by auto
qed

corollary *cond-exp-le-c:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$
*dered-real-vector}\}
assumes *integrable M f*
and *AE x in M. f x ≤ c*
shows *AE x in M. cond-exp M F f x ≤ c*
proof –
have *AE x in M. cond-exp M F f x = – cond-exp M F (λx. – f x) x* **using**
cond-exp-uminus[OF assms(1)] **by force**
moreover have *AE x in M. cond-exp M F (λx. – f x) x ≥ – c* **using** *assms*
by (*intro cond-exp-ge-c*) *auto*
ultimately show *?thesis* **by** (*force simp add: minus-le-iff*)
qed*

corollary *cond-exp-mono:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$
*dered-real-vector}\}
assumes *integrable M f integrable M g AE x in M. f x ≤ g x*
shows *AE x in M. cond-exp M F f x ≤ cond-exp M F g x*
using *cond-exp-le-c[OF Bochner-Integration.integrable-diff, OF assms(1,2), of*
0]
cond-exp-diff[OF assms(1,2)] assms(3) **by auto***

corollary *cond-exp-min:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$
*dered-real-vector}\}
assumes *integrable M f integrable M g*
shows *AE ξ in M. cond-exp M F (λx. min (f x) (g x)) ξ ≤ min (cond-exp M F*
f ξ) (cond-exp M F g ξ)
proof –
have *AE ξ in M. cond-exp M F (λx. min (f x) (g x)) ξ ≤ cond-exp M F f ξ* **by**
(intro cond-exp-mono integrable-min assms, simp)
moreover have *AE ξ in M. cond-exp M F (λx. min (f x) (g x)) ξ ≤ cond-exp*
M F g ξ **by** (*intro cond-exp-mono integrable-min assms, simp*)
ultimately show *AE ξ in M. cond-exp M F (λx. min (f x) (g x)) ξ ≤ min*
(cond-exp M F f ξ) (cond-exp M F g ξ) **by fastforce**
qed*

corollary *cond-exp-max:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$

dered-real-vector }
assumes *integrable M f integrable M g*
shows *AE ξ in M. cond-exp M F ($\lambda x. \max (f x) (g x)$) $\xi \geq \max (cond-exp M F f \xi) (cond-exp M F g \xi)$*
proof –
have *AE ξ in M. cond-exp M F ($\lambda x. \max (f x) (g x)$) $\xi \geq cond-exp M F f \xi$ by (intro cond-exp-mono integrable-max assms, simp)*
moreover have *AE ξ in M. cond-exp M F ($\lambda x. \max (f x) (g x)$) $\xi \geq cond-exp M F g \xi$ by (intro cond-exp-mono integrable-max assms, simp)*
ultimately show *AE ξ in M. cond-exp M F ($\lambda x. \max (f x) (g x)$) $\xi \geq \max (cond-exp M F f \xi) (cond-exp M F g \xi)$ by fastforce*
qed

corollary *cond-exp-inf*:

fixes *f :: 'a \Rightarrow 'b :: {second-countable-topology, banach, linorder-topology, ordered-real-vector, lattice}*
assumes *integrable M f integrable M g*
shows *AE ξ in M. cond-exp M F ($\lambda x. \inf (f x) (g x)$) $\xi \leq \inf (cond-exp M F f \xi) (cond-exp M F g \xi)$*
unfolding *inf-min using assms by (rule cond-exp-min)*

corollary *cond-exp-sup*:

fixes *f :: 'a \Rightarrow 'b :: {second-countable-topology, banach, linorder-topology, ordered-real-vector, lattice}*
assumes *integrable M f integrable M g*
shows *AE ξ in M. cond-exp M F ($\lambda x. \sup (f x) (g x)$) $\xi \geq \sup (cond-exp M F f \xi) (cond-exp M F g \xi)$*
unfolding *sup-max using assms by (rule cond-exp-max)*

end

end

theory *Stochastic-Process*

imports *Filtered-Measure Measure-Space-Addendum*

begin

0.8 Stochastic Process

A stochastic process is a collection of random variables, indexed by a type 'b.

locale *stochastic-process* =

fixes *M t₀ and X :: 'b :: {second-countable-topology, linorder-topology} \Rightarrow 'a \Rightarrow 'c :: {second-countable-topology, banach}*
assumes *random-variable[measurable]: $\bigwedge i. t_0 \leq i \implies X i \in \text{borel-measurable } M$*
begin

definition *left-continuous* **where** *left-continuous = (AE ξ in M. $\forall i. \text{continuous (at-left } i) (\lambda i. X i \xi)$)*

definition *right-continuous* **where** *right-continuous = (AE ξ in M. $\forall i. \text{continuous$*

(*at-right* i) ($\lambda i. X\ i\ \xi$)

end

locale *nat-stochastic-process* = *stochastic-process* $M\ 0 :: \text{nat } X$ **for** $M\ X$
locale *real-stochastic-process* = *stochastic-process* $M\ 0 :: \text{real } X$ **for** $M\ X$

context *stochastic-process*
begin

lemma *compose*:
assumes $\bigwedge i. t_0 \leq i \implies f\ i \in \text{borel-measurable borel}$
shows *stochastic-process* $M\ t_0\ (\lambda i\ \xi. (f\ i)\ (X\ i\ \xi))$
by (*unfold-locale*) (*intro measurable-compose* [*OF* *random-variable* *assms*])

lemma *norm*: *stochastic-process* $M\ t_0\ (\lambda i\ \xi. \text{norm}\ (X\ i\ \xi))$ **by** (*fastforce* *intro*:
compose)

lemma *scaleR-right*:
assumes *stochastic-process* $M\ t_0\ Y$
shows *stochastic-process* $M\ t_0\ (\lambda i\ \xi. (Y\ i\ \xi) *_{\mathbb{R}} (X\ i\ \xi))$
using *stochastic-process.random-variable* [*OF* *assms*] *random-variable* **by** (*unfold-locale*)
simp

lemma *scaleR-right-const-fun*:
assumes $f \in \text{borel-measurable } M$
shows *stochastic-process* $M\ t_0\ (\lambda i\ \xi. f\ \xi *_{\mathbb{R}} (X\ i\ \xi))$
by (*unfold-locale*) (*intro borel-measurable-scaleR* *assms* *random-variable*)

lemma *scaleR-right-const*: *stochastic-process* $M\ t_0\ (\lambda i\ \xi. c\ i *_{\mathbb{R}} (X\ i\ \xi))$
by (*unfold-locale*) *simp*

lemma *add*:
assumes *stochastic-process* $M\ t_0\ Y$
shows *stochastic-process* $M\ t_0\ (\lambda i\ \xi. X\ i\ \xi + Y\ i\ \xi)$
using *stochastic-process.random-variable* [*OF* *assms*] *random-variable* **by** (*unfold-locale*)
simp

lemma *diff*:
assumes *stochastic-process* $M\ t_0\ Y$
shows *stochastic-process* $M\ t_0\ (\lambda i\ \xi. X\ i\ \xi - Y\ i\ \xi)$
using *stochastic-process.random-variable* [*OF* *assms*] *random-variable* **by** (*unfold-locale*)
simp

lemma *uminus*: *stochastic-process* $M\ t_0\ (-X)$ **using** *scaleR-right-const* [*of* $\lambda -. -1$]
by (*simp* *add*: *fun-Compl-def*)

lemma *partial-sum*: *stochastic-process* $M\ t_0\ (\lambda n\ \xi. \sum_{i \in \{t_0..<n\}} X\ i\ \xi)$ **by** (*unfold-locale*)

simp

lemma *partial-sum'*: *stochastic-process* $M\ t_0\ (\lambda n\ \xi.\ \sum_{i \in \{t_0..n\}} X\ i\ \xi)$ **by** (*unfold-locales*) *simp*

end

lemma *stochastic-process-const-fun*:
assumes $f \in \text{borel-measurable } M$
shows *stochastic-process* $M\ t_0\ (\lambda i.\ f)$ **using** *assms* **by** (*unfold-locales*)

lemma *stochastic-process-const*:
shows *stochastic-process* $M\ t_0\ (\lambda i.\ c\ i)$ **by** (*unfold-locales*) *simp*

lemma *stochastic-process-sum*:
assumes $\bigwedge i.\ i \in I \implies \text{stochastic-process } M\ t_0\ (X\ i)$
shows *stochastic-process* $M\ t_0\ (\lambda k\ \xi.\ \sum_{i \in I} X\ i\ k\ \xi)$ **using** *assms* [*THEN stochastic-process.random-variable*] **by** (*unfold-locales, auto*)

0.8.1 Natural Filtration

The natural filtration induced by a stochastic process X is the filtration generated by all events involving the process up to the time index t , i.e. $\Sigma_t = \sigma\{X\ s \mid s \leq t\}$.

definition *natural-filtration* :: $'a\ \text{measure} \Rightarrow 'b \Rightarrow ('b \Rightarrow 'a \Rightarrow 'c :: \text{topological-space}) \Rightarrow 'b :: \{\text{second-countable-topology, linorder-topology}\} \Rightarrow 'a\ \text{measure}$
where

natural-filtration $M\ t_0\ Y = (\lambda t.\ \text{sigma-gen } (\text{space } M)\ \text{borel } \{Y\ i \mid i \in \{t_0..t\}\})$

context *stochastic-process*

begin

lemma *sets-natural-filtration'*: *sets* (*natural-filtration* $M\ t_0\ X\ t$) = *sigma-sets* (*space* M) $(\bigcup_{i \in \{t_0..t\}} \{X\ i - 'A \cap \text{space } M \mid A.\ A \in \text{borel}\})$
unfolding *natural-filtration-def* *sets-sigma-gen* **by** (*intro sigma-sets-eqI*) *blast+*

lemma

shows *sets-natural-filtration*: *sets* (*natural-filtration* $M\ t_0\ X\ t$) = *sigma-sets* (*space* M) $(\bigcup_{i \in \{t_0..t\}} \{X\ i - 'A \cap \text{space } M \mid A.\ \text{open } A\})$
and *space-natural-filtration*[*simp*]: *space* (*natural-filtration* $M\ t_0\ X\ t$) = *space* M

proof –

show *space* (*natural-filtration* $M\ t_0\ X\ t$) = *space* M **unfolding** *natural-filtration-def* *space-sigma-gen* ..

show *sets* (*natural-filtration* $M\ t_0\ X\ t$) = *sigma-sets* (*space* M) $(\bigcup_{i \in \{t_0..t\}} \{X\ i - 'A \cap \text{space } M \mid A.\ \text{open } A\})$ **unfolding** *sets-natural-filtration'*

proof (*intro sigma-sets-eqI, clarify*)

fix i **and** $A :: 'c\ \text{set}$ **assume** *asm*: $i \in \{t_0..t\}\ A \in \text{sets borel}$

hence $A \in \text{sigma-sets UNIV } \{S.\ \text{open } S\}$ **unfolding** *borel-def* **by** *simp*

thus $X \text{ } i \text{ } - ' A \cap \text{space } M \in \text{sigma-sets } (\text{space } M) (\bigcup i \in \{t_0..t\}. \{X \text{ } i \text{ } - ' A \cap \text{space } M \mid A. \text{open } A\})$
proof (*induction*)
case (*Compl a*)
have $X \text{ } i \text{ } - ' (UNIV - a) \cap \text{space } M = \text{space } M - (X \text{ } i \text{ } - ' a \cap \text{space } M)$ **by** *blast*
then show *?case using Compl(2)[THEN sigma-sets.Compl]* **by** *presburger*
next
case (*Union a*)
have $X \text{ } i \text{ } - ' \bigcup (\text{range } a) \cap \text{space } M = \bigcup (\text{range } (\lambda j. X \text{ } i \text{ } - ' a \text{ } j \cap \text{space } M))$
by *blast*
then show *?case using Union(2)[THEN sigma-sets.Union]* **by** *presburger*
qed (*auto intro: asm*)
qed (*intro sigma-sets.Basic, fastforce*)
qed

lemma *subalgebra-natural-filtration:*

shows *subalgebra M (natural-filtration M t₀ X i)*
unfolding *subalgebra-def using measurable-family-iff-contains-sigma-gen* **by** (*force simp add: natural-filtration-def*)

end

sublocale *stochastic-process \subseteq filtered-measure-natural-filtration: filtered-measure M natural-filtration M t₀ X t₀*

by (*unfold-locales*) (*intro subalgebra-natural-filtration, simp only: sets-natural-filtration, intro sigma-sets-subseteq, force*)

In order to show that the natural filtration constitutes a filtered sigma finite measure, we need to provide a countable exhausting set in the preimage of $X \text{ } t_0$.

lemma (*in sigma-finite-measure*) *sigma-finite-filtered-measure-natural-filtration:*

assumes *stochastic-process M t₀ X*
and *exhausting-set: countable A ($\bigcup A = \text{space } M \wedge a. a \in A \implies \text{emeasure } M \text{ } a \neq \infty \wedge a. a \in A \implies \exists b \in \text{borel}. a = X \text{ } t_0 - ' b \cap \text{space } M$)*

shows *sigma-finite-filtered-measure M (natural-filtration M t₀ X) t₀*

proof (*unfold-locales*)

interpret *stochastic-process M t₀ X* **by** (*rule assms*)

have $A \subseteq \text{sets } (\text{restr-to-subalg } M (\text{natural-filtration } M \text{ } t_0 \text{ } X \text{ } t_0))$ **using** *exhausting-set* **by** (*simp add: sets-restr-to-subalg[OF subalgebra-natural-filtration] sets-natural-filtration'*)
fast

moreover have $\bigcup A = \text{space } (\text{restr-to-subalg } M (\text{natural-filtration } M \text{ } t_0 \text{ } X \text{ } t_0))$

unfolding *space-restr-to-subalg* **using** *exhausting-set* **by** *simp*

moreover have $\forall a \in A. \text{emeasure } (\text{restr-to-subalg } M (\text{natural-filtration } M \text{ } t_0 \text{ } X \text{ } t_0)) \text{ } a \neq \infty$ **using** *calculation(1) exhausting-set(3)*

by (*auto simp add: sets-restr-to-subalg[OF subalgebra-natural-filtration] emea-sure-restr-to-subalg[OF subalgebra-natural-filtration]*)

ultimately show $\exists A. \text{countable } A \wedge A \subseteq \text{sets } (\text{restr-to-subalg } M (\text{natural-filtration } M \text{ } t_0 \text{ } X \text{ } t_0)) \wedge \bigcup A = \text{space } (\text{restr-to-subalg } M (\text{natural-filtration } M \text{ } t_0 \text{ } X \text{ } t_0)) \wedge$

$(\forall a \in A. \text{emeasure } (\text{restr-to-subalg } M \text{ (natural-filtration } M \ t_0 \ X \ t_0)) \ a \neq \infty)$ **using**
exhausting-set **by** *blast*
show $\bigwedge i j. \llbracket t_0 \leq i; i \leq j \rrbracket \implies \text{sets } (\text{natural-filtration } M \ t_0 \ X \ i) \subseteq \text{sets } (\text{natural-filtration } M \ t_0 \ X \ j)$ **using** *filtered-measure-natural-filtration.subalgebra-F* **by** (*simp add: sub-algebra-def*)
qed (*auto intro: stochastic-process.subalgebra-natural-filtration assms(1)*)

lemma (*in finite-measure*) *sigma-finite-filtered-measure-natural-filtration*:
assumes *stochastic-process* $M \ t_0 \ X$
shows *sigma-finite-filtered-measure* $M \text{ (natural-filtration } M \ t_0 \ X) \ t_0$
proof (*intro sigma-finite-filtered-measure-natural-filtration[OF assms(1), of {space M}]*)
have $\text{space } M = X \ t_0 - ' UNIV \cap \text{space } M$ **by** *blast*
thus $\bigwedge a. a \in \{\text{space } M\} \implies \exists b \in \text{sets borel}. a = X \ t_0 - ' b \cap \text{space } M$ **by** *force*
qed (*auto*)

0.9 Adapted Process

We call a collection a stochastic process X adapted if $X \ i$ is $F \ i$ -borel-measurable for all indices i .

locale *adapted-process* = *filtered-measure* $M \ F \ t_0$ **for** $M \ F \ t_0$ **and** $X :: - \Rightarrow - \Rightarrow -$
 $:: \{\text{second-countable-topology, banach}\} +$
assumes *adapted[measurable]*: $\bigwedge i. t_0 \leq i \implies X \ i \in \text{borel-measurable } (F \ i)$
begin

lemma *adaptedE[elim]*:
assumes $\llbracket \bigwedge j \ i. t_0 \leq j \implies j \leq i \implies X \ j \in \text{borel-measurable } (F \ i) \rrbracket \implies P$
shows P
using *assms* **using** *adapted* **by** (*metis dual-order.trans borel-measurable-subalgebra sets-F-mono space-F*)

lemma *adaptedD*:
assumes $t_0 \leq j \leq i$
shows $X \ j \in \text{borel-measurable } (F \ i)$ **using** *assms adaptedE* **by** *meson*

end

locale *nat-adapted-process* = *adapted-process* $M \ F \ 0 :: \text{nat } X$ **for** $M \ F \ X$
sublocale *nat-adapted-process* \subseteq *nat-filtered-measure* ..

locale *real-adapted-process* = *adapted-process* $M \ F \ 0 :: \text{real } X$ **for** $M \ F \ X$
sublocale *real-adapted-process* \subseteq *real-filtered-measure* ..

lemma (*in filtered-measure*) *adapted-process-const-fun*:
assumes $f \in \text{borel-measurable } (F \ t_0)$
shows *adapted-process* $M \ F \ t_0 \ (\lambda \cdot. f)$
using *measurable-from-subalg subalgebra-F assms* **by** (*unfold-locales*) *blast*

lemma (*in filtered-measure*) *adapted-process-const*:

shows *adapted-process* $M F t_0 (\lambda i -. c i)$ **by** (*unfold-locales*) *simp*

context *adapted-process*
begin

lemma *compose*:
assumes $\bigwedge i. f i \in \text{borel-measurable borel}$
shows *adapted-process* $M F t_0 (\lambda i \xi. (f i) (X i \xi))$
by (*unfold-locales*) (*intro measurable-compose*[*OF adapted assms*])

lemma *norm*: *adapted-process* $M F t_0 (\lambda i \xi. \text{norm } (X i \xi))$ **by** (*fastforce intro: compose*)

lemma *scaleR-right*:
assumes *adapted-process* $M F t_0 R$
shows *adapted-process* $M F t_0 (\lambda i \xi. (R i \xi) *_R (X i \xi))$
using *adapted-process.adapted*[*OF assms*] *adapted* **by** (*unfold-locales*) *simp*

lemma *scaleR-right-const-fun*:
assumes $f \in \text{borel-measurable } (F t_0)$
shows *adapted-process* $M F t_0 (\lambda i \xi. f \xi *_R (X i \xi))$
using *assms* **by** (*fast intro: scaleR-right adapted-process-const-fun*)

lemma *scaleR-right-const*: *adapted-process* $M F t_0 (\lambda i \xi. c i *_R (X i \xi))$ **by** (*unfold-locales*) *simp*

lemma *add*:
assumes *adapted-process* $M F t_0 Y$
shows *adapted-process* $M F t_0 (\lambda i \xi. X i \xi + Y i \xi)$
using *adapted-process.adapted*[*OF assms*] *adapted* **by** (*unfold-locales*) *simp*

lemma *diff*:
assumes *adapted-process* $M F t_0 Y$
shows *adapted-process* $M F t_0 (\lambda i \xi. X i \xi - Y i \xi)$
using *adapted-process.adapted*[*OF assms*] *adapted* **by** (*unfold-locales*) *simp*

lemma *uminus*: *adapted-process* $M F t_0 (-X)$ **using** *scaleR-right-const*[*of* $\lambda -. -1$]
by (*simp add: fun-Compl-def*)

lemma *partial-sum*: *adapted-process* $M F t_0 (\lambda n \xi. \sum_{i \in \{t_0..<n\}} X i \xi)$
proof (*unfold-locales*)
fix $i :: 'b$
have $X j \in \text{borel-measurable } (F i)$ **if** $t_0 \leq j < i$ **for** j **using** *that adaptedE* **by** *fastforce*
thus $(\lambda \xi. \sum_{i \in \{t_0..<i\}} X i \xi) \in \text{borel-measurable } (F i)$ **by** *simp*
qed

lemma *partial-sum'*: *adapted-process* $M F t_0 (\lambda n \xi. \sum_{i \in \{t_0..n\}} X i \xi)$
proof (*unfold-locales*)

```

fix  $i :: 'b$ 
have  $X j \in \text{borel-measurable } (F i) \text{ if } t_0 \leq j \text{ for } j \text{ using that adaptedE by}$ 
 $\text{meson}$ 
thus  $(\lambda \xi. \sum_{i \in \{t_0..i\}}. X i \xi) \in \text{borel-measurable } (F i) \text{ by simp}$ 
qed

end

lemma (in filtered-measure) adapted-process-sum:
assumes  $\bigwedge i. i \in I \implies \text{adapted-process } M F t_0 (X i)$ 
shows  $\text{adapted-process } M F t_0 (\lambda k \xi. \sum_{i \in I}. X i k \xi)$ 
proof -
{
fix  $i k$  assume  $i \in I$  and  $\text{asm: } t_0 \leq k$ 
then interpret  $\text{adapted-process } M F t_0 X i$  using assms by simp
have  $X i k \in \text{borel-measurable } M X i k \in \text{borel-measurable } (F k) \text{ using mea-}$ 
 $\text{surable-from-subalg subalgebra adapted asm by (blast, simp)}$ 
}
thus ?thesis by (unfold-locales) simp
qed

```

An adapted process is necessarily a stochastic process.

sublocale *adapted-process* \subseteq *stochastic-process* **using** *measurable-from-subalg sub-*
algebra adapted **by** (*unfold-locales*) *blast*

sublocale *nat-adapted-process* \subseteq *nat-stochastic-process* ..
sublocale *real-adapted-process* \subseteq *real-stochastic-process* ..

A stochastic process is always adapted to the natural filtration it generates.

sublocale *stochastic-process* \subseteq *adapted-natural: adapted-process* $M \text{ natural-filtration}$
 $M t_0 X t_0 X$ **by** (*unfold-locales*) (*auto simp add: natural-filtration-def intro: ran-*
dom-variable measurable-sigma-gen)

0.10 Progressively Measurable Process

locale *progressive-process* = *filtered-measure* $M F t_0$ **for** $M F t_0$ **and** $X :: - \Rightarrow -$
 $\Rightarrow - :: \{\text{second-countable-topology, banach}\} +$
assumes *progressive[measurable]*: $\bigwedge t. t_0 \leq t \implies (\lambda(i, x). X (\min t i) x) \in$
 $\text{borel-measurable } (\text{restrict-space borel } \{t_0..t\} \otimes_M F t)$
begin

lemma *progressiveD*:
assumes $S \in \text{borel}$
shows $(\lambda(j, \xi). X (\min i j) \xi) - ' S \cap (\{t_0..i\} \times \text{space } M) \in (\text{restrict-space borel}$
 $\{t_0..i\} \otimes_M F i)$
using *measurable-sets[OF progressive, OF - assms, of i]*
by (*cases* $t_0 \leq i$) (*auto simp add: space-F space-restrict-space sets-pair-measure*
space-pair-measure)

end

locale *nat-progressive-process* = *progressive-process* $M\ F\ 0 :: \text{nat } X$ **for** $M\ F\ X$
locale *real-progressive-process* = *progressive-process* $M\ F\ 0 :: \text{real } X$ **for** $M\ F\ X$

lemma (in *filtered-measure*) *prog-measurable-process-const-fun*:

assumes $f \in \text{borel-measurable } (F\ t_0)$
shows *progressive-process* $M\ F\ t_0\ (\lambda\cdot. f)$
proof (*unfold-locales*)
fix i **assume** *asm*: $t_0 \leq i$
have $f \in \text{borel-measurable } (F\ i)$ **using** *borel-measurable-mono*[*OF order.refl asm*]
assms **by** *blast*
thus *case-prod* $(\lambda\cdot. f) \in \text{borel-measurable } (\text{restrict-space borel } \{t_0..i\} \otimes_M F\ i)$
using *measurable-compose*[*OF measurable-snd*] **by** *simp*
qed

lemma (in *filtered-measure*) *prog-measurable-process-const*:

assumes $c \in \text{borel-measurable borel}$
shows *progressive-process* $M\ F\ t_0\ (\lambda i\ \cdot. c\ i)$
using *assms* **by** (*unfold-locales*) (*auto simp add: measurable-split-conv intro!: measurable-compose*[*OF measurable-fst*] *measurable-restrict-space1*)

context *progressive-process*

begin

lemma *compose*:

assumes *case-prod* $f \in \text{borel-measurable borel}$
shows *progressive-process* $M\ F\ t_0\ (\lambda i\ \xi. (f\ i)\ (X\ i\ \xi))$
proof
fix i **assume** *asm*: $t_0 \leq i$
have $(\lambda(j :: 'b, \xi :: 'a). (j, X\ (\min\ i\ j)\ \xi)) \in (\text{restrict-space borel } \{t_0..i\} \otimes_M F\ i) \rightarrow_M \text{borel} \otimes_M \text{borel}$ **using** *progressive*[*OF asm*] *measurable-fst'*[*OF measurable-restrict-space1, OF measurable-id*] **by** (*auto simp add: measurable-pair-iff space-pair-measure sets-pair-measure sets-restrict-space measurable-split-conv*)
moreover **have** $(\lambda(j :: 'b, y :: 'c). ((\min\ i\ j), y)) \in \text{borel} \otimes_M \text{borel} \rightarrow_M \text{borel} \otimes_M \text{borel}$ **by** *simp*
moreover **have** $(\lambda(j, \xi). f\ (\min\ i\ j)\ (X\ (\min\ i\ j)\ \xi)) = \text{case-prod } f\ o\ ((\lambda(j, y). ((\min\ i\ j), y))\ o\ (\lambda(j, \xi). (j, X\ (\min\ i\ j)\ \xi)))$ **by** *fastforce*
ultimately show $(\lambda(j :: 'b, \xi :: 'a). f\ (\min\ i\ j)\ (X\ (\min\ i\ j)\ \xi)) \in \text{borel-measurable } (\text{restrict-space borel } \{t_0..i\} \otimes_M F\ i)$ **unfolding** *borel-prod* **using** *assms measurable-comp*[*OF measurable-comp*] **by** *simp*
qed

lemma *norm*: *progressive-process* $M\ F\ t_0\ (\lambda i\ \xi. \text{norm } (X\ i\ \xi))$ **using** *measurable-compose*[*OF progressive borel-measurable-norm*] **by** (*unfold-locales*) *simp*

lemma *scaleR-right*:

assumes *progressive-process* $M\ F\ t_0\ R$
shows *progressive-process* $M\ F\ t_0\ (\lambda i\ \xi. (R\ i\ \xi) *_R (X\ i\ \xi))$

using *progressive-process.progressive*[*OF assms*] **by** (*unfold-locales*) (*simp add: progressive assms*)

lemma *scaleR-right-const-fun*:

assumes $f \in \text{borel-measurable } (F \ t_0)$

shows *progressive-process* $M \ F \ t_0 \ (\lambda i \ \xi. f \ \xi \ *_R \ (X \ i \ \xi))$

using *assms* **by** (*fast intro: scaleR-right prog-measurable-process-const-fun*)

lemma *scaleR-right-const*:

assumes $c \in \text{borel-measurable borel}$

shows *progressive-process* $M \ F \ t_0 \ (\lambda i \ \xi. c \ i \ *_R \ (X \ i \ \xi))$

using *assms* **by** (*fastforce intro: scaleR-right prog-measurable-process-const*)

lemma *add*:

assumes *progressive-process* $M \ F \ t_0 \ Y$

shows *progressive-process* $M \ F \ t_0 \ (\lambda i \ \xi. X \ i \ \xi + Y \ i \ \xi)$

using *progressive-process.progressive*[*OF assms*] **by** (*unfold-locales*) (*simp add: progressive assms*)

lemma *diff*:

assumes *progressive-process* $M \ F \ t_0 \ Y$

shows *progressive-process* $M \ F \ t_0 \ (\lambda i \ \xi. X \ i \ \xi - Y \ i \ \xi)$

using *progressive-process.progressive*[*OF assms*] **by** (*unfold-locales*) (*simp add: progressive assms*)

lemma *uminus: progressive-process* $M \ F \ t_0 \ (-X)$ **using** *scaleR-right-const*[*of* $\lambda-. -1$] **by** (*simp add: fun-Compl-def*)

end

A progressively measurable process is also adapted.

sublocale *progressive-process* \subseteq *adapted-process* **using** *measurable-compose-rev*[*OF progressive measurable-Pair1*] **unfolding** *prod.case* **by** (*unfold-locales*) *simp*

sublocale *nat-progressive-process* \subseteq *nat-adapted-process* ..

sublocale *real-progressive-process* \subseteq *real-adapted-process* ..

In the discrete setting, adaptedness is equivalent to progressive measurability.

sublocale *nat-adapted-process* \subseteq *nat-progressive-process*

proof (*unfold-locales*)

fix $i :: \text{nat}$

have $(\lambda(j, y). \min i \ j) - ' S \cap \{0..i\} \times \text{space } (F \ i) = S \times \text{space } (F \ i)$ **if** $S \in \text{restrict-space borel } \{0..i\}$ **for** S **using** *sets.sets-into-space*[*OF that*] **by** *auto*

hence $(\lambda(j, y). \min i \ j) \in \text{restrict-space borel } \{0..i\} \otimes_M F \ i \rightarrow_M \text{restrict-space borel } \{0..i\}$ **by** (*intro measurableI*) (*auto simp add: space-pair-measure sets-pair-measure*)

hence $(\lambda(j, x). (\min i \ j, x)) \in \text{restrict-space borel } \{0..i\} \otimes_M F \ i \rightarrow_M \text{restrict-space borel } \{0..i\} \otimes_M F \ i$ **by** (*intro measurable-pair*) *simp+*

moreover have $\text{case-prod } X \in \text{borel-measurable } (\text{restrict-space borel } \{0..i\} \otimes_M F i)$
proof (*intro borel-measurableI*)
fix $S :: 'b \text{ set}$ **assume** $\text{open-}S$: $\text{open } S$
{
fix j **assume** $\text{asm}: j \leq i$
hence $X j - ' S \cap \text{space } M \in F i$ **using** $\text{adaptedD}[of j, \text{THEN measurable-sets}]$
 $\text{space-}F \text{ open-}S$ **by** *fastforce*
moreover have $\text{case-prod } X - ' S \cap \{j\} \times \text{space } M = \{j\} \times (X j - ' S \cap \text{space } M)$ **for** j **by** *fast*
moreover have $\{j :: \text{nat}\} \in \text{restrict-space borel } \{0..i\}$ **using** asm **by** (*simp add: sets-restrict-space-iff*)
ultimately have $\text{case-prod } X - ' S \cap \{j\} \times \text{space } M \in \text{restrict-space borel } \{0..i\} \otimes_M F i$ **by** *simp*
}
hence $(\lambda j. (\lambda(x, y). X x y) - ' S \cap \{j\} \times \text{space } M) - ' \{0..i\} \subseteq \text{restrict-space borel } \{0..i\} \otimes_M F i$ **by** *blast*
moreover have $\text{case-prod } X - ' S \cap \text{space } (\text{restrict-space borel } \{0..i\} \otimes_M F i) = (\bigcup_{j \leq i}. \text{case-prod } X - ' S \cap \{j\} \times \text{space } M)$ **unfolding** $\text{space-pair-measure space-restrict-space space-}F$ **by** *force*
ultimately show $\text{case-prod } X - ' S \cap \text{space } (\text{restrict-space borel } \{0..i\} \otimes_M F i) \in \text{restrict-space borel } \{0..i\} \otimes_M F i$ **by** (*metis sets.countable-UN*)
qed
moreover have $(\lambda(x, y). X x y) o (\lambda(j, x). (\min i j, x)) = (\lambda(j, x). X (\min i j, x))$ **by** *fastforce*
ultimately show $(\lambda(j, x). X (\min i j, x)) \in \text{borel-measurable } (\text{restrict-space borel } \{0..i\} \otimes_M F i)$ **by** (*metis measurable-comp*)
qed

0.11 Predictable Process

We introduce the constant Σ_P to denote the predictable sigma algebra.

context *filtered-measure*
begin

definition $\Sigma_P :: ('b \times 'a) \text{ measure}$ **where** $\text{predictable-sigma}: \Sigma_P \equiv \text{sigma } (\{t_0..\} \times \text{space } M) (\{\{s <..t\} \times A \mid A \text{ s t. } A \in F s \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in F t_0\})$

lemma $\text{space-predictable-sigma}[simp]: \text{space } \Sigma_P = (\{t_0..\} \times \text{space } M)$ **unfolding** $\text{predictable-sigma space-measure-of-conv}$ **by** *blast*

lemma $\text{sets-predictable-sigma}: \text{sets } \Sigma_P = \text{sigma-sets } (\{t_0..\} \times \text{space } M) (\{\{s <..t\} \times A \mid A \text{ s t. } A \in F s \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in F t_0\})$

unfolding predictable-sigma **using** $\text{space-}F \text{ sets.sets-into-space}$ **by** (*subst sets-measure-of fastforce+*)

lemma $\text{measurable-predictable-sigma-snd}:$

assumes $\text{countable } \mathcal{I} \mathcal{I} \subseteq \{\{s <..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\} \{t_0 <..\} \subseteq (\bigcup \mathcal{I})$

shows $snd \in \Sigma_P \rightarrow_M F t_0$
proof (*intro measurableI, force simp add: space-F*)
fix $S :: 'a \text{ set}$ **assume** $asm: S \in F t_0$
have *countable*: *countable* $((\lambda I. I \times S) \text{ ' } \mathcal{I})$ **using** *assms(1)* **by** *blast*
have $(\lambda I. I \times S) \text{ ' } \mathcal{I} \subseteq \{\{s <..t\} \times A \mid A \text{ s } t. A \in F s \wedge t_0 \leq s \wedge s < t\}$ **using**
sets-F-mono[OF order-refl, THEN subsetD, OF - asm] *assms(2)* **by** *blast*
hence $(\bigcup I \in \mathcal{I}. I \times S) \cup \{t_0\} \times S \in \Sigma_P$ **unfolding** *sets-predictable-sigma* **using**
asm **by** (*intro sigma-sets-Un[OF sigma-sets-UNION[OF countable] sigma-sets.Basic]*
sigma-sets.Basic) *blast+*
moreover **have** $snd - ' S \cap \text{space } \Sigma_P = \{t_0..\} \times S$ **using** *sets.sets-into-space[OF*
asm] **by** (*fastforce simp add: space-F*)
moreover **have** $(\bigcup I \in \mathcal{I}. I \times S) \cup \{t_0\} \times S = \{t_0..\} \times S$ **using** *assms(2,3)*
using *ivl-disj-un(1)* **by** *fastforce*
ultimately show $snd - ' S \cap \text{space } \Sigma_P \in \Sigma_P$ **by** *argo*
qed

lemma *measurable-predictable-sigma-fst*:

assumes *countable* \mathcal{I} $\mathcal{I} \subseteq \{\{s <..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\}$ $\{t_0 <..\} \subseteq (\bigcup \mathcal{I})$
shows *fst* $\in \text{borel-measurable } \Sigma_P$
proof –
have $A \times \text{space } M \in \text{sets } \Sigma_P$ **if** $A \in \text{sigma-sets } \{t_0..\}$ $\{\{s <..t\} \mid s \text{ t. } t_0 \leq s \wedge s$
 $< t\}$ **for** A **unfolding** *sets-predictable-sigma* **using** *that*
proof (*induction rule: sigma-sets.induct*)
case (*Basic a*)
thus *?case* **using** *space-F sets.top* **by** *blast*
next
case (*Compl a*)
have $(\{t_0..\} - a) \times \text{space } M = \{t_0..\} \times \text{space } M - a \times \text{space } M$ **by** *blast*
then show *?case* **using** *Compl(2)[THEN sigma-sets.Compl]* **by** *presburger*
next
case (*Union a*)
have $\bigcup (\text{range } a) \times \text{space } M = \bigcup (\text{range } (\lambda i. a \ i \times \text{space } M))$ **by** *blast*
then show *?case* **using** *Union(2)[THEN sigma-sets.Union]* **by** *presburger*
qed (*auto*)
moreover **have** *restrict-space borel* $\{t_0..\} = \text{sigma } \{t_0..\} \{\{s <..t\} \mid s \text{ t. } t_0 \leq s$
 $\wedge s < t\}$
proof –
have *sigma-sets* $\{t_0..\} ((\cap) \{t_0..\} \text{ ' sigma-sets UNIV (range greaterThan)}) =$
sigma-sets $\{t_0..\} \{\{s <..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\}$
proof (*intro sigma-sets-eqI ; clarify*)
fix $A :: 'b \text{ set}$ **assume** $asm: A \in \text{sigma-sets UNIV (range greaterThan)}$
thus $\{t_0..\} \cap A \in \text{sigma-sets } \{t_0..\} \{\{s <..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\}$
proof (*induction rule: sigma-sets.induct*)
case (*Basic a*)
then obtain s **where** $a = \{s <..\}$ **by** *blast*
show *?case*
proof (*cases* $t_0 \leq s$)
case *True*
hence $*$: $\{t_0..\} \cap a = (\bigcup i \in \mathcal{I}. \{s <..\} \cap i)$ **using** *s assms(3)* **by** *force*

```

    have  $((\cap) \{s<..\} \text{ ' } \mathcal{I}) \subseteq \text{sigma-sets } \{t_0..\} \{\{s<..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\}$ 
    proof (clarify)
      fix A assume  $A \in \mathcal{I}$ 
      then obtain  $s' t'$  where  $A: A = \{s'<..t'\} \text{ } t_0 \leq s' s' < t'$  using assms(2)
    by blast
      hence  $\{s<..\} \cap A = \{\max s s' <..t'\}$  by fastforce
      moreover have  $t_0 \leq \max s s'$  using A True by linarith
      moreover have  $\max s s' < t'$  if  $s < t'$  using A that by linarith
      moreover have  $\{s<..\} \cap A = \{\}$  if  $\neg s < t'$  using A that by force
      ultimately show  $\{s<..\} \cap A \in \text{sigma-sets } \{t_0..\} \{\{s<..t\} \mid s \text{ t. } t_0 \leq s$ 
 $\wedge s < t\}$  by (cases  $s < t'$ ) (blast, simp add: sigma-sets.Empty)
      qed
      thus ?thesis unfolding * using assms(1) by (intro sigma-sets-UNION)
    auto
  next
    case False
    hence  $\{t_0..\} \cap a = \{t_0..\}$  using s by force
    thus ?thesis using sigma-sets-top by auto
  qed
next
case (Compl a)
have  $\{t_0..\} \cap (\text{UNIV} - a) = \{t_0..\} - (\{t_0..\} \cap a)$  by blast
then show ?case using Compl(2)[THEN sigma-sets.Compl] by presburger
next
case (Union a)
have  $\{t_0..\} \cap \bigcup (\text{range } a) = \bigcup (\text{range } (\lambda i. \{t_0..\} \cap a \ i))$  by blast
then show ?case using Union(2)[THEN sigma-sets.Union] by presburger
qed (simp add: sigma-sets.Empty)
next
fix s t assume asm:  $t_0 \leq s < t$ 
hence *:  $\{s<..t\} = \{s<..\} \cap (\{t_0..\} - \{t<..\})$  by force
have  $\{s<..\} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\} \text{ ' } \text{sigma-sets UNIV (range greaterThan)})$ 
using asm by (intro sigma-sets.Basic) auto
moreover have  $\{t_0..\} - \{t<..\} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\} \text{ ' } \text{sigma-sets
 $\text{UNIV (range greaterThan)})$  using asm by (intro sigma-sets.Compl sigma-sets.Basic)
auto
ultimately show  $\{s<..t\} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\} \text{ ' } \text{sigma-sets UNIV
 $(\text{range greaterThan}))$  unfolding * Int-range-binary[of  $\{s<..\}$ ] by (intro sigma-sets-Inter[OF
 $\text{- binary-in-sigma-sets}]$ ) auto
  qed
  thus ?thesis unfolding borel-Ioi restrict-space-def emeasure-sigma by (force
intro: sigma-eqI)
  qed
  ultimately have  $\text{restrict-space borel } \{t_0..\} \otimes_M \text{sigma (space } M) \{\} \subseteq \text{sets } \Sigma_P$ 

  unfolding sets-pair-measure space-restrict-space space-measure-of-conv
  using space-predictable-sigma sets.sigma-algebra-axioms[of  $\Sigma_P$ ]
  by (intro sigma-algebra.sigma-sets-subset) (auto simp add: sigma-sets-empty-eq
sets-measure-of-conv)$$ 
```

moreover have $\text{space } (\text{restrict-space borel } \{t_0..\} \otimes_M \text{sigma } (\text{space } M) \{ \}) = \text{space } \Sigma_P$ **by** (*simp add: space-pair-measure*)

moreover have $\text{fst} \in \text{restrict-space borel } \{t_0..\} \otimes_M \text{sigma } (\text{space } M) \{ \} \rightarrow_M \text{borel}$ **by** (*fastforce intro: measurable-fst''[OF measurable-restrict-space1, of $\lambda x. x$]*)

ultimately show *?thesis* **by** (*meson borel-measurable-subalgebra*)
qed

end

locale *predictable-process* = *filtered-measure* $M F t_0$ **for** $M F t_0$ **and** $X :: - \Rightarrow -$
 $\Rightarrow - :: \{ \text{second-countable-topology, banach} \} +$
assumes *predictable: case-prod* $X \in \text{borel-measurable } \Sigma_P$
begin

lemmas *predictableD* = *measurable-sets*[*OF predictable, unfolded space-predictable-sigma*]

end

locale *nat-predictable-process* = *predictable-process* $M F 0 :: \text{nat } X$ **for** $M F X$
locale *real-predictable-process* = *predictable-process* $M F 0 :: \text{real } X$ **for** $M F X$

lemma (**in** *nat-filtered-measure*) *measurable-predictable-sigma-snd*:
shows $\text{snd} \in \Sigma_P \rightarrow_M F 0$
by (*intro measurable-predictable-sigma-snd[of range $(\lambda x. \{ \text{Suc } x \})$]*) (*force | simp add: greaterThan-0*)**+**

lemma (**in** *real-filtered-measure*) *measurable-predictable-sigma-snd*:
shows $\text{snd} \in \Sigma_P \rightarrow_M F 0$
using *real-arch-simple* **by** (*intro measurable-predictable-sigma-snd[of range $(\lambda x :: \text{nat. } \{ 0 < .. \text{real } (\text{Suc } x) \})$]*) (*fastforce intro: add-increasing*)**+**

lemma (**in** *nat-filtered-measure*) *measurable-predictable-sigma-fst*:
shows $\text{fst} \in \text{borel-measurable } \Sigma_P$
by (*intro measurable-predictable-sigma-fst[of range $(\lambda x. \{ \text{Suc } x \})$]*) (*force | simp add: greaterThan-0*)**+**

lemma (**in** *real-filtered-measure*) *measurable-predictable-sigma-fst*:
shows $\text{fst} \in \text{borel-measurable } \Sigma_P$
using *real-arch-simple* **by** (*intro measurable-predictable-sigma-fst[of range $(\lambda x :: \text{nat. } \{ 0 < .. \text{real } (\text{Suc } x) \})$]*) (*fastforce intro: add-increasing*)**+**

lemma (**in** *filtered-measure*) *predictable-process-const-fun*:
assumes $\text{snd} \in \Sigma_P \rightarrow_M F t_0$ $f \in \text{borel-measurable } (F t_0)$
shows *predictable-process* $M F t_0$ $(\lambda \cdot. f)$

using *measurable-compose-rev*[*OF assms*(2)] *assms*(1) **by** (*unfold-locales*) (*auto simp add: measurable-split-conv*)

lemma (*in nat-filtered-measure*) *predictable-process-const-fun*:
assumes $f \in \text{borel-measurable } (F \ 0)$
shows *nat-predictable-process* $M \ F \ (\lambda \cdot. f)$
using *assms* **by** (*intro predictable-process-const-fun*[*OF measurable-predictable-sigma-snd*,
THEN nat-predictable-process.intro])

lemma (*in real-filtered-measure*) *predictable-process-const-fun*:
assumes $f \in \text{borel-measurable } (F \ 0)$
shows *real-predictable-process* $M \ F \ (\lambda \cdot. f)$
using *assms* **by** (*intro predictable-process-const-fun*[*OF measurable-predictable-sigma-snd*,
THEN real-predictable-process.intro])

lemma (*in filtered-measure*) *predictable-process-const*:
assumes $\text{fst} \in \text{borel-measurable } \Sigma_P \ c \in \text{borel-measurable borel}$
shows *predictable-process* $M \ F \ t_0 \ (\lambda i \cdot. c \ i)$
using *assms* **by** (*unfold-locales*) (*simp add: measurable-split-conv*)

lemma (*in filtered-measure*) *predictable-process-const'*:
shows *predictable-process* $M \ F \ t_0 \ (\lambda \cdot. \cdot. c)$
by (*unfold-locales*) *simp*

lemma (*in nat-filtered-measure*) *predictable-process-const*:
assumes $c \in \text{borel-measurable borel}$
shows *nat-predictable-process* $M \ F \ (\lambda i \cdot. c \ i)$
using *assms* **by** (*intro predictable-process-const*[*OF measurable-predictable-sigma-fst*,
THEN nat-predictable-process.intro])

lemma (*in real-filtered-measure*) *predictable-process-const*:
assumes $c \in \text{borel-measurable borel}$
shows *real-predictable-process* $M \ F \ (\lambda i \cdot. c \ i)$
using *assms* **by** (*intro predictable-process-const*[*OF measurable-predictable-sigma-fst*,
THEN real-predictable-process.intro])

context *predictable-process*
begin

lemma *compose*:
assumes $\text{fst} \in \text{borel-measurable } \Sigma_P \ \text{case-prod } f \in \text{borel-measurable borel}$
shows *predictable-process* $M \ F \ t_0 \ (\lambda i \ \xi. (f \ i) (X \ i \ \xi))$
proof
have $(\lambda(i, \xi). (i, X \ i \ \xi)) \in \Sigma_P \rightarrow_M \text{borel} \otimes_M \text{borel}$ **using** *predictable assms*(1)
by (*auto simp add: measurable-pair-iff measurable-split-conv*)
moreover **have** $(\lambda(i, \xi). f \ i (X \ i \ \xi)) = \text{case-prod } f \ o \ (\lambda(i, \xi). (i, X \ i \ \xi))$ **by**
fastforce
ultimately show $(\lambda(i, \xi). f \ i (X \ i \ \xi)) \in \text{borel-measurable } \Sigma_P$ **unfolding** *borel-prod*
using *assms* **by** *simp*

qed

lemma *norm*: *predictable-process* $M F t_0 (\lambda i \xi. \text{norm } (X i \xi))$ **using** *measurable-compose*[*OF predictable borel-measurable-norm*]
by (*unfold-locales*) (*simp add: prod.case-distrib*)

lemma *scaleR-right*:
assumes *predictable-process* $M F t_0 R$
shows *predictable-process* $M F t_0 (\lambda i \xi. (R i \xi) *_{\mathbb{R}} (X i \xi))$
using *predictable predictable-process.predictable*[*OF assms*] **by** (*unfold-locales*)
(*auto simp add: measurable-split-conv*)

lemma *scaleR-right-const-fun*:
assumes $\text{snd} \in \Sigma_P \rightarrow_M F t_0 f \in \text{borel-measurable } (F t_0)$
shows *predictable-process* $M F t_0 (\lambda i \xi. f \xi *_{\mathbb{R}} (X i \xi))$
using *assms* **by** (*fast intro: scaleR-right predictable-process-const-fun*)

lemma *scaleR-right-const*:
assumes $\text{fst} \in \text{borel-measurable } \Sigma_P c \in \text{borel-measurable borel}$
shows *predictable-process* $M F t_0 (\lambda i \xi. c i *_{\mathbb{R}} (X i \xi))$
using *assms* **by** (*fastforce intro: scaleR-right predictable-process-const*)

lemma *scaleR-right-const'*: *predictable-process* $M F t_0 (\lambda i \xi. c *_{\mathbb{R}} (X i \xi))$
by (*fastforce intro: scaleR-right predictable-process-const'*)

lemma *add*:
assumes *predictable-process* $M F t_0 Y$
shows *predictable-process* $M F t_0 (\lambda i \xi. X i \xi + Y i \xi)$
using *predictable predictable-process.predictable*[*OF assms*] **by** (*unfold-locales*)
(*auto simp add: measurable-split-conv*)

lemma *diff*:
assumes *predictable-process* $M F t_0 Y$
shows *predictable-process* $M F t_0 (\lambda i \xi. X i \xi - Y i \xi)$
using *predictable predictable-process.predictable*[*OF assms*] **by** (*unfold-locales*)
(*auto simp add: measurable-split-conv*)

lemma *uminus*: *predictable-process* $M F t_0 (-X)$ **using** *scaleR-right-const'*[*of -1*]
by (*simp add: fun-Compl-def*)

end

Every predictable process is also progressively measurable.

sublocale *predictable-process* \subseteq *progressive-process*

proof (*unfold-locales*)

fix $i :: 'b$ **assume** *asm*: $t_0 \leq i$

let $?min = (\lambda(j, x). (\min i j, x))$

{

fix $S :: ('b \times 'a) \text{ set}$ **assume** $S \in \{\{s <..t\} \times A \mid A s t. A \in F s \wedge t_0 \leq s \wedge s$

$< t\} \cup \{\{t_0\} \times A \mid A. A \in F t_0\}$
hence $?min - ' S \cap (\{t_0..i\} \times space\ M) \in restrict\text{-}space\ borel\ \{t_0..i\} \otimes_M F\ i$
proof
assume $S \in \{\{s<..t\} \times A \mid A\ s\ t. A \in F\ s \wedge t_0 \leq s \wedge s < t\}$
then obtain $s\ t\ A$ **where** $S\text{-}is: S = \{s<..t\} \times A\ t_0 \leq s\ s < t\ A \in F\ s$ **by**
blast
hence $?min - ' S \cap (\{t_0..i\} \times space\ M) = \{s<..min\ i\ t\} \times A$ **using**
sets.sets-into-space[OF S-is(4)] by (auto simp add: space-F)
then show $?thesis$ **using** $S\text{-}is\ sets\text{-}F\text{-}mono[of\ s\ i]$ **by** $(cases\ s \leq i)\ (fastforce\ simp\ add: sets\text{-}restrict\text{-}space\text{-}iff)+$
next
assume $S \in \{\{t_0\} \times A \mid A. A \in F\ t_0\}$
then obtain A **where** $S\text{-}is: S = \{t_0\} \times A\ A \in F\ t_0$ **by** *blast*
hence $?min - ' S \cap (\{t_0..i\} \times space\ M) = \{t_0\} \times A$ **using** *asm sets.sets-into-space[OF S-is(2)] by (auto simp add: space-F)*
thus $?thesis$ **using** $S\text{-}is\ sets\text{-}F\text{-}mono[OF\ order\text{-}refl\ asm]$ *asm* **by** $(fastforce\ simp\ add: sets\text{-}restrict\text{-}space\text{-}iff)$
qed
hence $?min - ' S \cap space\ (restrict\text{-}space\ borel\ \{t_0..i\} \otimes_M F\ i) \in restrict\text{-}space\ borel\ \{t_0..i\} \otimes_M F\ i$ **by** $(simp\ add: space\text{-}pair\text{-}measure\ space\text{-}F[OF\ asm])$
}
moreover have $\{\{s<..t\} \times A \mid A\ s\ t. A \in sets\ (F\ s) \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in sets\ (F\ t_0)\} \subseteq Pow\ (\{t_0..i\} \times space\ M)$ **using** *sets.sets-into-space by (fastforce simp add: space-F)*
ultimately have $?min \in restrict\text{-}space\ borel\ \{t_0..i\} \otimes_M F\ i \rightarrow_M \Sigma_P$ **using** *space-F[OF asm] by (intro measurable-sigma-sets[OF sets-predictable-sigma]) (fast, force simp add: space-pair-measure)*
moreover have $case\text{-}prod\ X\ o\ ?min = (\lambda(j, x). X\ (min\ i\ j)\ x)$ **by** *fastforce*
ultimately show $case\text{-}prod\ (\lambda j. X\ (min\ i\ j)) \in borel\text{-}measurable\ (restrict\text{-}space\ borel\ \{t_0..i\} \otimes_M F\ i)$ **by** $(metis\ measurable\text{-}comp\ predictable)$
qed

sublocale $nat\text{-}predictable\text{-}process \subseteq nat\text{-}progressive\text{-}process ..$
sublocale $real\text{-}predictable\text{-}process \subseteq real\text{-}progressive\text{-}process ..$

0.12 Additional Lemmas for Discrete Time Processes

lemma **(in** $nat\text{-}adapted\text{-}process$ $)$ $partial\text{-}sum\text{-}Suc: nat\text{-}adapted\text{-}process\ M\ F\ (\lambda n\ \xi. \sum_{i < n}. X\ (Suc\ i)\ \xi)$
proof $(unfold\ locales)$
fix i
have $X\ j \in borel\text{-}measurable\ (F\ i)$ **if** $j \leq i$ **for** j **using** *that adaptedD* **by** *blast*
thus $(\lambda \xi. \sum_{i < i}. X\ (Suc\ i)\ \xi) \in borel\text{-}measurable\ (F\ i)$ **by** *auto*
qed

The following lemma characterizes predictability in a discrete-time setting.

lemma **(in** $nat\text{-}filtered\text{-}measure$ $)$ $sets\text{-}in\text{-}filtration:$
assumes $(\bigcup i. \{i\} \times A\ i) \in \Sigma_P$
shows $A\ (Suc\ i) \in F\ i\ A\ 0 \in F\ 0$

```

using assms unfolding sets-predictable-sigma
proof (induction  $(\bigcup i. \{i\} \times A \ i)$  arbitrary: A)
  case Basic
  {
    assume  $\exists S. (\bigcup i. \{i\} \times A \ i) = \{0\} \times S$ 
    then obtain S where S:  $(\bigcup i. \{i\} \times A \ i) = \{bot\} \times S$  unfolding bot-nat-def
  by blast
    hence  $S \in F \ bot$  using Basic by (fastforce simp add: times-eq-iff bot-nat-def)
    moreover have  $A \ i = \{\}$  if  $i \neq bot$  for i using that S by blast
    moreover have  $A \ bot = S$  using S by blast
    ultimately have  $A \ (Suc \ i) \in F \ i \ A \ 0 \in F \ 0$  for i unfolding bot-nat-def by
    (auto simp add: bot-nat-def)
  }
  note  $*$  = this
  {
    assume  $\nexists S. (\bigcup i. \{i\} \times A \ i) = \{0\} \times S$ 
    then obtain s t B where B:  $(\bigcup i. \{i\} \times A \ i) = \{s<..t\} \times B \ B \in sets \ (F \ s)$ 
  s < t using Basic by auto
    hence  $A \ i = B$  if  $i \in \{s<..t\}$  for i using that by fast
    moreover have  $A \ i = \{\}$  if  $i \notin \{s<..t\}$  for i using B that by fastforce
    ultimately have  $A \ (Suc \ i) \in F \ i \ A \ 0 \in F \ 0$  for i unfolding bot-nat-def using
B sets-F-mono by (auto simp add: bot-nat-def) (metis less-Suc-eq-le sets.empty-sets
subset-eq)
  }
  note  $**$  = this
  show  $A \ (Suc \ i) \in sets \ (F \ i) \ A \ 0 \in sets \ (F \ 0)$  using  $*(1)[of \ i] \ *(2) \ **(1)[of \ i]$ 
 $**(2)$  by blast+
next
  case Empty
  {
    case 1
    then show ?case using Empty by simp
  next
    case 2
    then show ?case using Empty by simp
  }
next
  case (Compl a)
  have a-in:  $a \subseteq \{0..\} \times space \ M$  using Compl(1) sets.sets-into-space sets-predictable-sigma
space-predictable-sigma by metis
  hence A-in:  $A \ i \subseteq space \ M$  for i using Compl(4) by blast
  have a:  $a = \{0..\} \times space \ M - (\bigcup i. \{i\} \times A \ i)$  using a-in Compl(4) by blast
  also have  $... = (\bigcup j. \{j\} \times (space \ M - A \ j))$  by blast
  finally have  $*$ :  $(space \ M - A \ (Suc \ i)) \in F \ i \ (space \ M - A \ 0) \in F \ 0$  using
Compl(2,3) by auto
  {
    case 1
    then show ?case using  $*$  A-in by (metis bot-nat-0.extremum double-diff
sets.Diff sets.top sets-F-mono sets-le-imp-space-le space-F)
  }

```



```

next
  case 2
    then show ?case using * A-in by (metis bot-nat-0.extremum double-diff
sets.Diff sets.top sets-F-mono sets-le-imp-space-le space-F)
  }
next
  case (Union a)
    have a-in:  $a \cap i \subseteq \{0..\} \times \text{space } M$  for  $i$  using Union(1) sets.sets-into-space
sets-predictable-sigma space-predictable-sigma by metis
    hence A-in:  $A \cap i \subseteq \text{space } M$  for  $i$  using Union(4) by blast
    have snd  $x \in \text{snd } ' (a \cap i \cap (\{fst\ x\} \times \text{space } M))$  if  $x \in a \cap i$  for  $i$   $x$  using that
a-in by fastforce
    hence a-i:  $a \cap i = (\bigcup j. \{j\} \times (\text{snd } ' (a \cap i \cap (\{j\} \times \text{space } M))))$  for  $i$  by force
    have A-i:  $A \cap i = \text{snd } ' (\bigcup (\text{range } a) \cap (\{i\} \times \text{space } M))$  for  $i$  unfolding
Union(4) using A-in by force
    have *:  $\text{snd } ' (a \cap j \cap (\{Suc\ i\} \times \text{space } M)) \in F\ i$   $\text{snd } ' (a \cap j \cap (\{0\} \times \text{space } M))$ 
 $\in F\ 0$  for  $j$  using Union(2,3)[OF a-i] by auto
    {
      case 1
        have  $(\bigcup j. \text{snd } ' (a \cap j \cap (\{Suc\ i\} \times \text{space } M))) \in F\ i$  using * by fast
        moreover have  $(\bigcup j. \text{snd } ' (a \cap j \cap (\{Suc\ i\} \times \text{space } M))) = \text{snd } ' (\bigcup (\text{range } a) \cap$ 
a)  $\cap (\{Suc\ i\} \times \text{space } M))$  by fast
        ultimately show ?case using A-i by metis
      next
        case 2
          have  $(\bigcup j. \text{snd } ' (a \cap j \cap (\{0\} \times \text{space } M))) \in F\ 0$  using * by fast
          moreover have  $(\bigcup j. \text{snd } ' (a \cap j \cap (\{0\} \times \text{space } M))) = \text{snd } ' (\bigcup (\text{range } a) \cap$ 
 $\{0\} \times \text{space } M))$  by fast
          ultimately show ?case using A-i by metis
        }
    }
qed

```

This leads to the following useful fact.

theorem (in *nat-predictable-process*) *adapted-Suc*: *nat-adapted-process* $M\ F\ (\lambda i.$
 $X\ (Suc\ i))$

proof (*unfold-locales, intro borel-measurableI*)

fix $S :: 'b\ set$ and i assume *open-S*: *open* S

have $\{Suc\ i\} = \{i < .. Suc\ i\}$ by fastforce

hence $\{Suc\ i\} \times \text{space } M \in \Sigma_P$ unfolding *space-F[symmetric, of i]* sets-predictable-sigma
by (*intro sigma-sets.Basic*) blast

moreover have *case-prod* $X - ' S \cap (UNIV \times \text{space } M) \in \Sigma_P$ unfolding
atLeast-0[symmetric] using *open-S* by (*intro predictableD, simp add: borel-open*)

ultimately have *case-prod* $X - ' S \cap (\{Suc\ i\} \times \text{space } M) \in \Sigma_P$ unfolding
sets-predictable-sigma using *space-F* sets.sets-into-space

by (*subst Times-Int-distrib1[of {Suc i} UNIV space M, simplified]*, *subst*
inf commute, subst Int-assoc[symmetric], *subst Int-range-binary*)

(*intro sigma-sets-Inter binary-in-sigma-sets, fast*) +

moreover have *case-prod* $X - ' S \cap (\{Suc\ i\} \times \text{space } M) = \{Suc\ i\} \times (X\ (Suc$
 $i) - ' S \cap \text{space } M)$ by (*auto simp add: le-Suc-eq*)

moreover have ... = $(\bigcup j. \{j\} \times (\text{if } j = \text{Suc } i \text{ then } (X (\text{Suc } i) - 'S \cap \text{space } M) \text{ else } \{\}))$ **by** (*auto split: if-splits*)
ultimately have $(\bigcup j. \{j\} \times (\text{if } j = \text{Suc } i \text{ then } (X (\text{Suc } i) - 'S \cap \text{space } M) \text{ else } \{\})) \in \Sigma_P$ **by** *argo*
thus $X (\text{Suc } i) - 'S \cap \text{space } (F i) \in \text{sets } (F i)$ **using** *sets-in-filtration*[of $\lambda j. \text{if } j = \text{Suc } i \text{ then } (X (\text{Suc } i) - 'S \cap \text{space } M) \text{ else } \{\}$] *space-F* **by** *presburger*
qed

0.13 Processes with an Ordering

These locales are useful in the definition of sub- and supermartingales.

locale *stochastic-process-order* = *stochastic-process* $M t_0 X$ **for** $M t_0$ **and** $X :: - \Rightarrow - \Rightarrow - :: \{\text{linorder-topology, ordered-real-vector}\}$
locale *adapted-process-order* = *adapted-process* $M F t_0 X$ **for** $M F t_0$ **and** $X :: - \Rightarrow - \Rightarrow - :: \{\text{linorder-topology, ordered-real-vector}\}$
locale *progressive-process-order* = *progressive-process* $M F t_0 X$ **for** $M F t_0$ **and** $X :: - \Rightarrow - \Rightarrow - :: \{\text{linorder-topology, ordered-real-vector}\}$
locale *predictable-process-order* = *predictable-process* $M F t_0 X$ **for** $M F t_0$ **and** $X :: - \Rightarrow - \Rightarrow - :: \{\text{linorder-topology, ordered-real-vector}\}$

locale *nat-stochastic-process-order* = *stochastic-process-order* $M 0 :: \text{nat } X$ **for** $M X$
locale *nat-adapted-process-order* = *adapted-process-order* $M F 0 :: \text{nat } X$ **for** $M F X$
locale *nat-progressive-process-order* = *progressive-process-order* $M F 0 :: \text{nat } X$ **for** $M F X$
locale *nat-predictable-process-order* = *predictable-process-order* $M F 0 :: \text{nat } X$ **for** $M F X$

locale *real-stochastic-process-order* = *stochastic-process-order* $M 0 :: \text{real } X$ **for** $M X$
locale *real-adapted-process-order* = *adapted-process-order* $M F 0 :: \text{real } X$ **for** $M F X$
locale *real-progressive-process-order* = *progressive-process-order* $M F 0 :: \text{real } X$ **for** $M F X$
locale *real-predictable-process-order* = *predictable-process-order* $M F 0 :: \text{real } X$ **for** $M F X$

sublocale *predictable-process-order* \subseteq *progressive-process-order* ..
sublocale *progressive-process-order* \subseteq *adapted-process-order* ..
sublocale *adapted-process-order* \subseteq *stochastic-process-order* ..

sublocale *nat-predictable-process-order* \subseteq *nat-progressive-process-order* ..

sublocale *nat-progressive-process-order* \subseteq *nat-adapted-process-order* ..
sublocale *nat-adapted-process-order* \subseteq *nat-stochastic-process-order* ..

sublocale *real-predictable-process-order* \subseteq *real-progressive-process-order* ..
sublocale *real-progressive-process-order* \subseteq *real-adapted-process-order* ..
sublocale *real-adapted-process-order* \subseteq *real-stochastic-process-order* ..

0.14 Processes with a Sigma Finite Filtration

locale *sigma-finite-adapted-process* = *adapted-process* + *sigma-finite-filtered-measure*
locale *sigma-finite-progressive-process* = *progressive-process* + *sigma-finite-filtered-measure*
locale *sigma-finite-predictable-process* = *predictable-process* + *sigma-finite-filtered-measure*

locale *sigma-finite-adapted-process-order* = *adapted-process-order* + *sigma-finite-filtered-measure*
locale *sigma-finite-progressive-process-order* = *progressive-process-order* + *sigma-finite-filtered-measure*
locale *sigma-finite-predictable-process-order* = *predictable-process-order* + *sigma-finite-filtered-measure*

locale *nat-sigma-finite-adapted-process* = *sigma-finite-adapted-process* *M F 0* :: *nat*
X for M F X

locale *nat-sigma-finite-progressive-process* = *sigma-finite-progressive-process* *M F*
0 :: *nat X for M F X*

locale *nat-sigma-finite-predictable-process* = *sigma-finite-predictable-process* *M F*
0 :: *nat X for M F X*

locale *nat-sigma-finite-adapted-process-order* = *sigma-finite-adapted-process-order*
M F 0 :: *nat X for M F X*

locale *nat-sigma-finite-progressive-process-order* = *sigma-finite-progressive-process-order*
M F 0 :: *nat X for M F X*

locale *nat-sigma-finite-predictable-process-order* = *sigma-finite-predictable-process-order*
M F 0 :: *nat X for M F X*

locale *real-sigma-finite-adapted-process* = *sigma-finite-adapted-process* *M F 0* ::
real X for M F X

locale *real-sigma-finite-progressive-process* = *sigma-finite-progressive-process* *M F*
0 :: *real X for M F X*

locale *real-sigma-finite-predictable-process* = *sigma-finite-predictable-process* *M F*
0 :: *real X for M F X*

locale *real-sigma-finite-adapted-process-order* = *sigma-finite-adapted-process-order*
M F 0 :: *real X for M F X*

locale *real-sigma-finite-progressive-process-order* = *sigma-finite-progressive-process-order*
M F 0 :: *real X for M F X*

locale *real-sigma-finite-predictable-process-order* = *sigma-finite-predictable-process-order*
M F 0 :: *real X for M F X*

```

sublocale sigma-finite-predictable-process  $\subseteq$  sigma-finite-progressive-process ..
sublocale sigma-finite-progressive-process  $\subseteq$  sigma-finite-adapted-process ..

sublocale nat-sigma-finite-predictable-process  $\subseteq$  nat-sigma-finite-progressive-process
..
sublocale nat-sigma-finite-progressive-process  $\subseteq$  nat-sigma-finite-adapted-process
..

sublocale real-sigma-finite-predictable-process  $\subseteq$  real-sigma-finite-progressive-process
..
sublocale real-sigma-finite-progressive-process  $\subseteq$  real-sigma-finite-adapted-process
..

sublocale nat-sigma-finite-adapted-process  $\subseteq$  nat-sigma-finite-filtered-measure ..
sublocale real-sigma-finite-adapted-process  $\subseteq$  real-sigma-finite-filtered-measure ..

sublocale sigma-finite-predictable-process-order  $\subseteq$  sigma-finite-progressive-process-order
..
sublocale sigma-finite-progressive-process-order  $\subseteq$  sigma-finite-adapted-process-order
..

sublocale nat-sigma-finite-predictable-process-order  $\subseteq$  nat-sigma-finite-progressive-process-order
..
sublocale nat-sigma-finite-progressive-process-order  $\subseteq$  nat-sigma-finite-adapted-process-order
..

sublocale real-sigma-finite-predictable-process-order  $\subseteq$  real-sigma-finite-progressive-process-order
..
sublocale real-sigma-finite-progressive-process-order  $\subseteq$  real-sigma-finite-adapted-process-order
..

sublocale nat-sigma-finite-adapted-process-order  $\subseteq$  nat-sigma-finite-filtered-measure
..
sublocale real-sigma-finite-adapted-process-order  $\subseteq$  real-sigma-finite-filtered-measure
..

Thus, right from the outset, we have pretty much every locale we may need.
end
theory Martingale
  imports Stochastic-Process Conditional-Expectation-Banach
begin

```

0.15 Martingale

```

locale martingale = sigma-finite-adapted-process +

```

assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
and *martingale-property*: $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE \xi \text{ in } M. X i \xi =$
cond-exp $M (F i) (X j) \xi$

locale *martingale-order* = *martingale* $M F t_0 X$ **for** $M F t_0$ **and** $X :: - \Rightarrow - \Rightarrow -$
 $:: \{ \text{linorder-topology, ordered-real-vector} \}$

lemma (**in** *sigma-finite-filtered-measure*) *martingale-const*[*intro*]:
assumes *integrable* $M f f \in \text{borel-measurable } (F t_0)$
shows *martingale* $M F t_0 (\lambda-. f)$
using *assms sigma-finite-subalgebra.cond-exp-F-meas*[*OF - assms*(1), *THEN AE-symmetric*]
borel-measurable-mono
by (*unfold-locale*) *blast*+

lemma (**in** *sigma-finite-filtered-measure*) *martingale-cond-exp*[*intro*]:
assumes *integrable* $M f$
shows *martingale* $M F t_0 (\lambda i. \text{cond-exp } M (F i) f)$
using *sigma-finite-subalgebra.borel-measurable-cond-exp'* *borel-measurable-cond-exp*

by (*unfold-locale*) (*auto intro: sigma-finite-subalgebra.cond-exp-nested-subalg*[*OF - assms*]
simp add: subalgebra-F subalgebra)

lemma (**in** *sigma-finite-filtered-measure*) *martingale-zero*[*intro*]: *martingale* $M F$
 $t_0 (\lambda-. 0)$ **by** *fastforce*

0.16 Submartingale

locale *submartingale* = *sigma-finite-adapted-process-order* +
assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
and *submartingale-property*: $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE \xi \text{ in } M. X i \xi \leq$
cond-exp $M (F i) (X j) \xi$

sublocale *martingale-order* \subseteq *submartingale* **using** *martingale-property* **by** (*unfold-locale*)
(*force simp add: integrable*)+

0.17 Supermartingale

locale *supermartingale* = *sigma-finite-adapted-process-order* +
assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
and *supermartingale-property*: $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE \xi \text{ in } M. X i \xi \geq$
cond-exp $M (F i) (X j) \xi$

sublocale *martingale-order* \subseteq *supermartingale* **using** *martingale-property* **by** (*unfold-locale*)
(*force simp add: integrable*)+

lemma *martingale-iff*: *martingale* $M F t_0 X \longleftrightarrow$ *submartingale* $M F t_0 X \wedge$
supermartingale $M F t_0 X$

proof (*rule iffI*)

assume *asm*: *martingale* $M F t_0 X$

interpret *martingale-order* $M F t_0 X$ **by** (*intro martingale-order.intro asm*)

```

  show submartingale  $M F t_0 X \wedge$  supermartingale  $M F t_0 X$  using submartin-
  gale-axioms supermartingale-axioms by blast
next
  assume asm: submartingale  $M F t_0 X \wedge$  supermartingale  $M F t_0 X$ 
  interpret submartingale  $M F t_0 X$  by (simp add: asm)
  interpret supermartingale  $M F t_0 X$  by (simp add: asm)
  show martingale  $M F t_0 X$  using submartingale-property supermartingale-property
  by (unfold-locales) (intro integrable, blast, force)
qed

```

0.18 Martingale Lemmas

```

context martingale
begin

```

```

lemma set-integral-eq:
  assumes  $A \in F i t_0 \leq i i \leq j$ 
  shows set-lebesgue-integral  $M A (X i) =$  set-lebesgue-integral  $M A (X j)$ 
proof -
  interpret sigma-finite-subalgebra  $M F i$  using assms(2) by blast
  have  $\int x \in A. X i x \partial M = \int x \in A. \text{cond-exp } M (F i) (X j) x \partial M$  using martin-
  gale-property[OF assms(2,3)] borel-measurable-cond-exp' assms subalgebra subalge-
  bra-def by (intro set-lebesgue-integral-cong-AE[OF - random-variable]) fastforce+
  also have  $\dots = \int x \in A. X j x \partial M$  using assms by (auto simp: integrable intro:
  cond-exp-set-integral[symmetric])
  finally show ?thesis .
qed

```

```

lemma scaleR-const[intro]:
  shows martingale  $M F t_0 (\lambda i x. c *_{\mathbb{R}} X i x)$ 
proof -
  {
    fix  $i j :: 'b$  assume asm:  $t_0 \leq i i \leq j$ 
    interpret sigma-finite-subalgebra  $M F i$  using asm by blast
    have  $\text{AE } x \text{ in } M. c *_{\mathbb{R}} X i x = \text{cond-exp } M (F i) (\lambda x. c *_{\mathbb{R}} X j x) x$  us-
    ing asm cond-exp-scaleR-right[OF integrable, of j, THEN AE-symmetric] martin-
    gale-property[OF asm] by force
  }
  thus ?thesis by (unfold-locales) (auto simp add: integrable martingale.integrable)
qed

```

```

lemma uminus[intro]:
  shows martingale  $M F t_0 (- X)$ 
  using scaleR-const[of -1] by (force intro: back-subst[of martingale  $M F t_0$ ])

```

```

lemma add[intro]:
  assumes martingale  $M F t_0 Y$ 
  shows martingale  $M F t_0 (\lambda i \xi. X i \xi + Y i \xi)$ 
proof -

```

```

interpret Y: martingale M F t0 Y by (rule assms)
{
  fix i j :: 'b assume asm: t0 ≤ i i ≤ j
  hence AE ξ in M. X i ξ + Y i ξ = cond-exp M (F i) (λx. X j x + Y j x) ξ
  using sigma-finite-subalgebra.cond-exp-add[OF - integrable martingale.integrable[OF
  assms], of F i j j, THEN AE-symmetric]
    martingale-property[OF asm] martingale.martingale-property[OF assms
  asm] by force
}
thus ?thesis using assms
by (unfold-locales) (auto simp add: integrable martingale.integrable)
qed

```

```

lemma diff[intro]:
  assumes martingale M F t0 Y
  shows martingale M F t0 (λi x. X i x - Y i x)
proof -
  interpret Y: martingale M F t0 Y by (rule assms)
  {
    fix i j :: 'b assume asm: t0 ≤ i i ≤ j
    hence AE ξ in M. X i ξ - Y i ξ = cond-exp M (F i) (λx. X j x - Y j x) ξ
    using sigma-finite-subalgebra.cond-exp-diff[OF - integrable martingale.integrable[OF
    assms], of F i j j, THEN AE-symmetric]
      martingale-property[OF asm] martingale.martingale-property[OF assms
    asm] by fastforce
  }
  thus ?thesis using assms by (unfold-locales) (auto simp add: integrable martin-
  gale.integrable)
qed

```

end

```

lemma (in sigma-finite-adapted-process) martingale-of-set-integral-eq:
  assumes integrable: ⋀i. integrable M (X i)
  and ⋀A i j. t0 ≤ i ⇒ i ≤ j ⇒ A ∈ F i ⇒ set-lebesgue-integral M A (X
  i) = set-lebesgue-integral M A (X j)
  shows martingale M F t0 X
proof (unfold-locales)
  fix i j :: 'b assume asm: t0 ≤ i i ≤ j
  interpret sigma-finite-subalgebra M F i using asm by blast
  interpret r: sigma-finite-measure restr-to-subalg M (F i) by (simp add: sigma-fin-subalg)
  {
    fix A assume A ∈ restr-to-subalg M (F i)
    hence *: A ∈ F i using sets-restr-to-subalg subalgebra asm by blast
    have set-lebesgue-integral (restr-to-subalg M (F i)) A (X i) = set-lebesgue-integral
    M A (X i) using * subalg asm by (auto simp: set-lebesgue-integral-def intro: inte-
    gral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator)
    also have ... = set-lebesgue-integral M A (cond-exp M (F i) (X j)) using *
    assms(2)[OF asm] cond-exp-set-integral[OF integrable] by auto
  }

```

finally have *set-lebesgue-integral* (*restr-to-subalg* M (F i)) A (X i) = *set-lebesgue-integral* (*restr-to-subalg* M (F i)) A (*cond-exp* M (F i) (X j)) **using** * *subalg* **by** (*auto simp*: *set-lebesgue-integral-def* *intro!*: *integral-subalgebra2*[*symmetric*] *borel-measurable-scaleR* *borel-measurable-cond-exp* *borel-measurable-indicator*)
}
hence $\text{AE } \xi \text{ in } \text{restr-to-subalg } M \text{ (} F \text{ } i \text{). } X \text{ } i \text{ } \xi = \text{cond-exp } M \text{ (} F \text{ } i \text{) (} X \text{ } j \text{)}$
 ξ **using** *asm* **by** (*intro* *r.density-unique*, *auto* *intro*: *integrable-in-subalg* *subalg* *borel-measurable-cond-exp* *integrable*)
thus $\text{AE } \xi \text{ in } M. X \text{ } i \text{ } \xi = \text{cond-exp } M \text{ (} F \text{ } i \text{) (} X \text{ } j \text{)}$ ξ **using** *AE-restr-to-subalg*[*OF* *subalg*] **by** *blast*
qed (*simp* *add*: *integrable*)

0.19 Submartingale Lemmas

context *submartingale*
begin

lemma *set-integral-le*:
assumes $A \in F$ i $t_0 \leq i$ $i \leq j$
shows *set-lebesgue-integral* M A (X i) \leq *set-lebesgue-integral* M A (X j)
using *submartingale-property*[*OF* *assms*(2), *of* j] *assms* *subalgebra*
by (*subst* *sigma-finite-subalgebra.cond-exp-set-integral*[*OF* - *integrable* *assms*(1), *of* j])
(*auto* *intro!*: *scaleR-left-mono* *integral-mono-AE-banach* *integrable-mult-indicator* *integrable* *simp* *add*: *subalgebra-def* *set-lebesgue-integral-def*)

lemma *cond-exp-diff-nonneg*:
assumes $t_0 \leq i$ $i \leq j$
shows $\text{AE } x \text{ in } M. 0 \leq \text{cond-exp } M \text{ (} F \text{ } i \text{) } (\lambda \xi. X \text{ } j \text{ } \xi - X \text{ } i \text{ } \xi) \text{ } x$
using *submartingale-property*[*OF* *assms*] *assms* *sigma-finite-subalgebra.cond-exp-diff*[*OF* - *integrable*(1,1), *of* - j i] *sigma-finite-subalgebra.cond-exp-F-meas*[*OF* - *integrable* *adapted*, *of* i] **by** *fastforce*

lemma *add*[*intro*]:
assumes *submartingale* M F t_0 Y
shows *submartingale* M F t_0 $(\lambda i \xi. X \text{ } i \text{ } \xi + Y \text{ } i \text{ } \xi)$
proof –
interpret Y : *submartingale* M F t_0 Y **by** (*rule* *assms*)
{
fix i j :: ' b **assume** *asm*: $t_0 \leq i$ $i \leq j$
hence $\text{AE } \xi \text{ in } M. X \text{ } i \text{ } \xi + Y \text{ } i \text{ } \xi \leq \text{cond-exp } M \text{ (} F \text{ } i \text{) } (\lambda x. X \text{ } j \text{ } x + Y \text{ } j \text{ } x) \text{ } \xi$
using *sigma-finite-subalgebra.cond-exp-add*[*OF* - *integrable* *submartingale.integrable*[*OF* *assms*], *of* F i j]
submartingale-property[*OF* *asm*] *submartingale.submartingale-property*[*OF* *assms* *asm*] *add-mono*[*of* $X \text{ } i$ - $Y \text{ } i$ -] **by** *force*
}
thus *?thesis* **using** *assms* **by** (*unfold-locales*) (*auto* *simp* *add*: *borel-measurable-add* *random-variable* *adapted* *integrable* Y .*random-variable* Y .*adapted* *submartingale.integrable*)

qed

lemma *diff[intro]*:

assumes *supermartingale* $M\ F\ t_0\ Y$

shows *submartingale* $M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi - Y\ i\ \xi)$

proof –

interpret Y : *supermartingale* $M\ F\ t_0\ Y$ **by** (rule *assms*)

{

fix $i\ j :: 'b$ **assume** *asm*: $t_0 \leq i\ i \leq j$

hence $AE\ \xi\ in\ M. X\ i\ \xi - Y\ i\ \xi \leq cond_exp\ M\ (F\ i)\ (\lambda x. X\ j\ x - Y\ j\ x)\ \xi$

using *sigma-finite-subalgebra.cond-exp-diff*[*OF* - *integrable supermartingale.integrable*[*OF assms*], *of F i j j*]

submartingale-property[*OF asm*] *supermartingale.supermartingale-property*[*OF assms asm*] *diff-mono*[*of X i - - Y i -*] **by** *force*

}

thus ?thesis **using** *assms* **by** (*unfold-locales*) (*auto simp add: borel-measurable-diff random-variable adapted integrable Y.random-variable Y.adapted supermartingale.integrable*)

qed

lemma *scaleR-nonneg*:

assumes $c \geq 0$

shows *submartingale* $M\ F\ t_0\ (\lambda i\ \xi. c *_{\mathbb{R}} X\ i\ \xi)$

proof

{

fix $i\ j :: 'b$ **assume** *asm*: $t_0 \leq i\ i \leq j$

thus $AE\ \xi\ in\ M. c *_{\mathbb{R}} X\ i\ \xi \leq cond_exp\ M\ (F\ i)\ (\lambda \xi. c *_{\mathbb{R}} X\ j\ \xi)\ \xi$

using *sigma-finite-subalgebra.cond-exp-scaleR-right*[*OF* - *integrable, of F i j c*] *submartingale-property*[*OF asm*] **by** (*fastforce intro!*: *scaleR-left-mono*[*OF - assms*])

}

qed (*auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable random-variable adapted borel-measurable-const-scaleR*)

lemma *scaleR-nonpos*:

assumes $c \leq 0$

shows *supermartingale* $M\ F\ t_0\ (\lambda i\ \xi. c *_{\mathbb{R}} X\ i\ \xi)$

proof

{

fix $i\ j :: 'b$ **assume** *asm*: $t_0 \leq i\ i \leq j$

thus $AE\ \xi\ in\ M. c *_{\mathbb{R}} X\ i\ \xi \geq cond_exp\ M\ (F\ i)\ (\lambda \xi. c *_{\mathbb{R}} X\ j\ \xi)\ \xi$

using *sigma-finite-subalgebra.cond-exp-scaleR-right*[*OF* - *integrable, of F i j c*] *submartingale-property*[*OF asm*]

by (*fastforce intro!*: *scaleR-left-mono-neg*[*OF - assms*])

}

qed (*auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable random-variable adapted borel-measurable-const-scaleR*)

lemma *uminus[intro]*:

shows *supermartingale* $M F t_0 (- X)$
unfolding *fun-Compl-def* **using** *scaleR-nonpos[of -1]* **by** *simp*

lemma *max*:
assumes *submartingale* $M F t_0 Y$
shows *submartingale* $M F t_0 (\lambda i \xi. \max (X i \xi) (Y i \xi))$
proof (*unfold-locales*)
interpret Y : *submartingale* $M F t_0 Y$ **by** (*rule assms*)
{
 fix $i j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
 have $AE \xi \text{ in } M. \max (X i \xi) (Y i \xi) \leq \max (\text{cond-exp } M (F i) (X j) \xi)$
 $(\text{cond-exp } M (F i) (Y j) \xi)$ **using** *submartingale-property* Y .*submartingale-property*
 asm **unfolding** *max-def* **by** *fastforce*
 thus $AE \xi \text{ in } M. \max (X i \xi) (Y i \xi) \leq \text{cond-exp } M (F i) (\lambda \xi. \max (X j \xi) (Y j \xi)) \xi$ **using** *sigma-finite-subalgebra.cond-exp-max[OF - integrable Y.integrable, of F i j j]* *asm* **by** (*fast intro: order.trans*)
}
show $\bigwedge i. t_0 \leq i \implies (\lambda \xi. \max (X i \xi) (Y i \xi)) \in \text{borel-measurable } (F i) \bigwedge i. t_0 \leq i \implies \text{integrable } M (\lambda \xi. \max (X i \xi) (Y i \xi))$ **by** (*force intro: Y.integrable integrable assms*)
qed

lemma *max-0*:
shows *submartingale* $M F t_0 (\lambda i \xi. \max 0 (X i \xi))$
proof -
interpret $zero$: *martingale-order* $M F t_0 \lambda -. 0$ **by** (*force intro: martingale-order.intro*)
show *?thesis* **by** (*intro zero.max submartingale-axioms*)
qed

end

lemma (*in sigma-finite-adapted-process-order*) *submartingale-of-cond-exp-diff-nonneg*:
assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
and *diff-nonneg*: $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE x \text{ in } M. 0 \leq \text{cond-exp } M (F i) (\lambda \xi. X j \xi - X i \xi) x$
shows *submartingale* $M F t_0 X$
proof (*unfold-locales*)
{
 fix $i j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
 thus $AE \xi \text{ in } M. X i \xi \leq \text{cond-exp } M (F i) (X j) \xi$
 using *diff-nonneg[OF asm]* *sigma-finite-subalgebra.cond-exp-diff[OF - integrable(1,1), of F i j i]*
 sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] **by** *fastforce*
}
qed (*intro integrable*)

lemma (*in sigma-finite-adapted-process-order*) *submartingale-of-set-integral-le*:
assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$

```

    and  $\bigwedge A \ i \ j. \ t_0 \leq i \implies i \leq j \implies A \in F \ i \implies \text{set-lebesgue-integral } M \ A \ (X \ i) \leq \text{set-lebesgue-integral } M \ A \ (X \ j)$ 
    shows submartingale  $M \ F \ t_0 \ X$ 
  proof (unfold-locales)
  {
    fix  $i \ j :: 'b$  assume asm:  $t_0 \leq i \ i \leq j$ 
    interpret  $r$ : sigma-finite-measure restr-to-subalg  $M \ (F \ i)$  using asm sigma-finite-subalgebra.sigma-fin-subalg
  by blast
  {
    fix  $A$  assume  $A \in \text{restr-to-subalg } M \ (F \ i)$ 
    hence *:  $A \in F \ i$  using asm sets-restr-to-subalg subalgebra by blast
    have  $\text{set-lebesgue-integral } (\text{restr-to-subalg } M \ (F \ i)) \ A \ (X \ i) = \text{set-lebesgue-integral } M \ A \ (X \ i)$  using * asm subalgebra by (auto simp: set-lebesgue-integral-def intro: integral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator)
    also have  $\dots \leq \text{set-lebesgue-integral } M \ A \ (\text{cond-exp } M \ (F \ i) \ (X \ j))$  using * assms(2)[OF asm] asm sigma-finite-subalgebra.cond-exp-set-integral[OF - integrable] by fastforce
    also have  $\dots = \text{set-lebesgue-integral } (\text{restr-to-subalg } M \ (F \ i)) \ A \ (\text{cond-exp } M \ (F \ i) \ (X \ j))$  using * asm subalgebra by (auto simp: set-lebesgue-integral-def intro!: integral-subalgebra2[symmetric] borel-measurable-scaleR borel-measurable-cond-exp borel-measurable-indicator)
    finally have  $0 \leq \text{set-lebesgue-integral } (\text{restr-to-subalg } M \ (F \ i)) \ A \ (\lambda \xi. \text{cond-exp } M \ (F \ i) \ (X \ j) \ \xi - X \ i \ \xi)$  using * asm subalgebra by (subst set-integral-diff, auto simp add: set-integrable-def sets-restr-to-subalg intro!: integrable adapted integrable-in-subalg borel-measurable-scaleR borel-measurable-indicator borel-measurable-cond-exp integrable-mult-indicator)
  }
  hence  $A \ E \ \xi$  in  $\text{restr-to-subalg } M \ (F \ i)$ .  $0 \leq \text{cond-exp } M \ (F \ i) \ (X \ j) \ \xi - X \ i \ \xi$ 
  by (intro r.density-nonneg integrable-in-subalg asm subalgebra borel-measurable-diff borel-measurable-cond-exp adapted Bochner-Integration.integrable-diff integrable-cond-exp integrable)
  thus  $A \ E \ \xi$  in  $M$ .  $X \ i \ \xi \leq \text{cond-exp } M \ (F \ i) \ (X \ j) \ \xi$  using  $A \ E \ \text{restr-to-subalg}[OF \text{subalgebra}] \text{asm}$  by simp
  }
qed (intro integrable)

```

0.20 Supermartingale Lemmas

The following lemmas are exact duals of submartingale lemmas.

context *supermartingale*
begin

lemma *set-integral-ge*:

```

  assumes  $A \in F \ i \ t_0 \leq i \ i \leq j$ 
  shows  $\text{set-lebesgue-integral } M \ A \ (X \ i) \geq \text{set-lebesgue-integral } M \ A \ (X \ j)$ 
  using supermartingale-property[OF assms(2), of  $j$ ] assms subalgebra
  by (subst sigma-finite-subalgebra.cond-exp-set-integral[OF - integrable assms(1), of  $j$ ])
  (auto intro!: scaleR-left-mono integral-mono-AE-banach integrable-mult-indicator)

```

integrable simp add: subalgebra-def set-lebesgue-integral-def)

lemma *cond-exp-diff-nonneg*:

assumes $t_0 \leq i \leq j$
shows $AE\ x\ in\ M. 0 \leq cond_exp\ M\ (F\ i)\ (\lambda \xi. X\ i\ \xi - X\ j\ \xi)\ x$
using *assms supermartingale-property*[*OF assms*] *sigma-finite-subalgebra.cond-exp-diff*[*OF*
- integrable(1,1), of F i i j]
sigma-finite-subalgebra.cond-exp-F-meas[*OF - integrable adapted, of i*] **by**
fastforce

lemma *add[intro]*:

assumes *supermartingale* $M\ F\ t_0\ Y$
shows *supermartingale* $M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi + Y\ i\ \xi)$
proof –
interpret Y : *supermartingale* $M\ F\ t_0\ Y$ **by** (*rule assms*)
{
fix $i\ j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
hence $AE\ \xi\ in\ M. X\ i\ \xi + Y\ i\ \xi \geq cond_exp\ M\ (F\ i)\ (\lambda x. X\ j\ x + Y\ j\ x)\ \xi$
using *sigma-finite-subalgebra.cond-exp-add*[*OF - integrable supermartingale.integrable*[*OF*
assms], *of F i j j*]
supermartingale-property[*OF asm*] *supermartingale.supermartingale-property*[*OF*
assms asm] *add-mono*[*of - X i - - Y i -*] **by** *force*
}
thus *?thesis using assms by* (*unfold-locals*) (*auto simp add: borel-measurable-add*
random-variable adapted integrable Y.random-variable Y.adapted supermartingale.integrable)

qed

lemma *diff[intro]*:

assumes *submartingale* $M\ F\ t_0\ Y$
shows *supermartingale* $M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi - Y\ i\ \xi)$
proof –
interpret Y : *submartingale* $M\ F\ t_0\ Y$ **by** (*rule assms*)
{
fix $i\ j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
hence $AE\ \xi\ in\ M. X\ i\ \xi - Y\ i\ \xi \geq cond_exp\ M\ (F\ i)\ (\lambda x. X\ j\ x - Y\ j\ x)\ \xi$
using *sigma-finite-subalgebra.cond-exp-diff*[*OF - integrable submartingale.integrable*[*OF*
assms], *of F i j j, unfolded fun-diff-def*]
supermartingale-property[*OF asm*] *submartingale.submartingale-property*[*OF*
assms asm] *diff-mono*[*of - X i - Y i -*] **by** *force*
}
thus *?thesis using assms by* (*unfold-locals*) (*auto simp add: borel-measurable-diff*
random-variable adapted integrable Y.random-variable Y.adapted submartingale.integrable)

qed

lemma *scaleR-nonneg*:

assumes $c \geq 0$
shows *supermartingale* $M\ F\ t_0\ (\lambda i\ \xi. c *_{\mathbb{R}} X\ i\ \xi)$

```

proof
{
  fix  $i\ j :: 'b$  assume  $asm: t_0 \leq i\ i \leq j$ 
  thus  $AE\ \xi\ in\ M. c *_R X\ i\ \xi \geq cond\_exp\ M\ (F\ i)\ (\lambda\xi. c *_R X\ j\ \xi)\ \xi$ 
  using  $\sigma\text{-finite-subalgebra.cond-exp-scaleR-right}[OF - integrable, of\ F\ i\ j\ c]$ 
   $supermartingale-property[OF\ asm]$  by  $(fastforce\ intro!: scaleR\text{-left-mono}[OF - assms])$ 
}
qed  $(auto\ simp\ add: borel\text{-measurable-integrable}\ borel\text{-measurable-scaleR}\ integrable\ random\text{-variable}\ adapted\ borel\text{-measurable-const-scaleR})$ 

```

lemma $scaleR\text{-nonpos}$:

```

assumes  $c \leq 0$ 
shows  $submartingale\ M\ F\ t_0\ (\lambda i\ \xi. c *_R X\ i\ \xi)$ 

```

```

proof
{
  fix  $i\ j :: 'b$  assume  $asm: t_0 \leq i\ i \leq j$ 
  thus  $AE\ \xi\ in\ M. c *_R X\ i\ \xi \leq cond\_exp\ M\ (F\ i)\ (\lambda\xi. c *_R X\ j\ \xi)\ \xi$ 
  using  $\sigma\text{-finite-subalgebra.cond-exp-scaleR-right}[OF - integrable, of\ F\ i\ j\ c]$ 
   $supermartingale-property[OF\ asm]$  by  $(fastforce\ intro!: scaleR\text{-left-mono-neg}[OF - assms])$ 
}
qed  $(auto\ simp\ add: borel\text{-measurable-integrable}\ borel\text{-measurable-scaleR}\ integrable\ random\text{-variable}\ adapted\ borel\text{-measurable-const-scaleR})$ 

```

lemma $uminus[intro]$:

```

shows  $submartingale\ M\ F\ t_0\ (-\ X)$ 
unfolding  $fun\text{-Compl-def}$  using  $scaleR\text{-nonpos}[of\ -1]$  by  $simp$ 

```

lemma min :

```

assumes  $supermartingale\ M\ F\ t_0\ Y$ 
shows  $supermartingale\ M\ F\ t_0\ (\lambda i\ \xi. min\ (X\ i\ \xi)\ (Y\ i\ \xi))$ 
proof  $(unfold\ locales)$ 
  interpret  $Y: supermartingale\ M\ F\ t_0\ Y$  by  $(rule\ assms)$ 
  {
    fix  $i\ j :: 'b$  assume  $asm: t_0 \leq i\ i \leq j$ 
    have  $AE\ \xi\ in\ M. min\ (X\ i\ \xi)\ (Y\ i\ \xi) \geq min\ (cond\_exp\ M\ (F\ i)\ (X\ j)\ \xi)\ (cond\_exp\ M\ (F\ i)\ (Y\ j)\ \xi)$ 
    using  $supermartingale-property\ Y.supermartingale-property\ asm$ 
    unfolding  $min\text{-def}$  by  $fastforce$ 
    thus  $AE\ \xi\ in\ M. min\ (X\ i\ \xi)\ (Y\ i\ \xi) \geq cond\_exp\ M\ (F\ i)\ (\lambda\xi. min\ (X\ j\ \xi)\ (Y\ j\ \xi))\ \xi$ 
    using  $\sigma\text{-finite-subalgebra.cond-exp-min}[OF - integrable\ Y.integrable, of\ F\ i\ j\ j]\ asm$ 
    by  $(fast\ intro: order.trans)$ 
  }
  show  $\bigwedge i. t_0 \leq i \implies (\lambda\xi. min\ (X\ i\ \xi)\ (Y\ i\ \xi)) \in borel\text{-measurable}\ (F\ i) \bigwedge i. t_0 \leq i \implies integrable\ M\ (\lambda\xi. min\ (X\ i\ \xi)\ (Y\ i\ \xi))$ 
  by  $(force\ intro: Y.integrable\ integrable\ assms)+$ 
qed

```

lemma $min-0$:

```

    shows supermartingale M F t0 (λ i ξ. min 0 (X i ξ))
  proof -
    interpret zero: martingale-order M F t0 λ- -. 0 by (force intro: martingale-order.intro)
    show ?thesis by (intro zero.min supermartingale-axioms)
  qed

end

```

lemma (in *sigma-finite-adapted-process-order*) *supermartingale-of-cond-exp-diff-nonneg*:

```

    assumes integrable: ∧ i. t0 ≤ i ⇒ integrable M (X i)
    and diff-nonneg: ∧ i j. t0 ≤ i ⇒ i ≤ j ⇒ AE x in M. 0 ≤ cond-exp M (F
i) (λ ξ. X i ξ - X j ξ) x
    shows supermartingale M F t0 X
  proof
    {
      fix i j :: 'b assume asm: t0 ≤ i i ≤ j
      thus AE ξ in M. X i ξ ≥ cond-exp M (F i) (X j) ξ
        using diff-nonneg[OF asm] sigma-finite-subalgebra.cond-exp-diff[OF - inte-
grable(1,1), of F i i j]
        sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] by
fastforce
    }
  qed (intro integrable)

```

lemma (in *sigma-finite-adapted-process-order*) *supermartingale-of-set-integral-ge*:

```

    assumes integrable: ∧ i. t0 ≤ i ⇒ integrable M (X i)
    and A i j. t0 ≤ i ⇒ i ≤ j ⇒ A ∈ F i ⇒ set-lebesgue-integral M A (X
j) ≤ set-lebesgue-integral M A (X i)
    shows supermartingale M F t0 X
  proof -
    interpret -: adapted-process M F t0 -X by (rule uminus)
    interpret uminus-X: sigma-finite-adapted-process-order M F t0 -X ..
    note * = set-integral-uminus[unfolded set-integrable-def, OF integrable-mult-indicator[OF
- integrable]]
    have supermartingale M F t0 (-(- X))
      using ord-eq-le-trans[OF * ord-le-eq-trans[OF le-imp-neg-le[OF assms(2)]] * [symmetric]]
subalgebra
    by (intro submartingale.uminus uminus-X.submartingale-of-set-integral-le)
      (clarsimp simp add: fun-Compl-def subalgebra-def integrable | fastforce)+
    thus ?thesis unfolding fun-Compl-def by simp
  qed

```

0.21 Discrete Time Martingales

```

locale nat-martingale = martingale M F 0 :: nat X for M F X
locale nat-submartingale = submartingale M F 0 :: nat X for M F X
locale nat-supermartingale = supermartingale M F 0 :: nat X for M F X

```

0.22 Discrete Time Martingales

lemma (in *nat-martingale*) *predictable-eq-zero*:

assumes *nat-predictable-process* $M F X$

shows $AE \xi \text{ in } M. X i \xi = X 0 \xi$

proof (*induction i*)

case 0

then show *?case* **by** (*simp add: bot-nat-def*)

next

case (*Suc i*)

interpret $S: \text{nat-adapted-process } M F \lambda i. X (Suc i)$ **by** (*intro nat-predictable-process.adapted-Suc assms*)

show *?case* **using** *Suc S.adapted[of i] martingale-property[OF - le-SucI, of i] sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable, of F i Suc i]* **by** *fastforce*
qed

lemma (in *nat-sigma-finite-adapted-process*) *martingale-of-set-integral-eq-Suc*:

assumes *integrable*: $\bigwedge i. \text{integrable } M (X i)$

and $\bigwedge A i. A \in F i \implies \text{set-lebesgue-integral } M A (X i) = \text{set-lebesgue-integral } M A (X (Suc i))$

shows *nat-martingale* $M F X$

proof (*intro nat-martingale.intro martingale-of-set-integral-eq*)

fix $i j A$ **assume** *asm*: $i \leq j \wedge A \in \text{sets } (F i)$

show $\text{set-lebesgue-integral } M A (X i) = \text{set-lebesgue-integral } M A (X j)$ **using** *asm*

proof (*induction j - i arbitrary: i j*)

case 0

then show *?case* **using** *asm* **by** *simp*

next

case (*Suc n*)

hence $*$: $n = j - Suc i$ **by** *linarith*

have $Suc i \leq j$ **using** *Suc(2,3)* **by** *linarith*

thus *?case* **using** *sets-F-mono[OF - le-SucI] Suc(4) Suc(1)[OF *]* **by** (*auto intro: assms(2)[THEN trans]*)

qed

qed (*simp add: integrable*)

lemma (in *nat-sigma-finite-adapted-process*) *martingale-nat*:

assumes *integrable*: $\bigwedge i. \text{integrable } M (X i)$

and $\bigwedge i. AE \xi \text{ in } M. X i \xi = \text{cond-exp } M (F i) (X (Suc i)) \xi$

shows *nat-martingale* $M F X$

proof (*unfold-locales*)

fix $i j :: \text{nat}$ **assume** *asm*: $i \leq j$

show $AE \xi \text{ in } M. X i \xi = \text{cond-exp } M (F i) (X j) \xi$ **using** *asm*

proof (*induction j - i arbitrary: i j*)

case 0

hence $j = i$ **by** *simp*

thus *?case* **using** *sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, THEN AE-symmetric]* **by** *blast*

next

case (*Suc n*)
have *j*: *j* = *Suc* (*n* + *i*) **using** *Suc* **by** *linarith*
have *n*: *n* = *n* + *i* - *i* **using** *Suc* **by** *linarith*
have *: *AE* ξ *in* *M*. *cond-exp* *M* (*F* (*n* + *i*)) (*X j*) ξ = *X* (*n* + *i*) ξ **unfolding**
j **using** *assms*(2)[*THEN AE-symmetric*] **by** *blast*
have *AE* ξ *in* *M*. *cond-exp* *M* (*F i*) (*X j*) ξ = *cond-exp* *M* (*F i*) (*cond-exp* *M*
(*F* (*n* + *i*)) (*X j*)) ξ **by** (*intro cond-exp-nested-subalg integrable subalg, simp add:*
subalgebra-def space-F sets-F-mono)
hence *AE* ξ *in* *M*. *cond-exp* *M* (*F i*) (*X j*) ξ = *cond-exp* *M* (*F i*) (*X* (*n* + *i*))
 ξ **using** *cond-exp-cong-AE*[*OF integrable-cond-exp integrable **] **by** *force*
thus ?*case* **using** *Suc*(1)[*OF n*] **by** *fastforce*
qed
qed (*simp add: integrable*)

lemma (*in nat-sigma-finite-adapted-process*) *martingale-of-cond-exp-diff-Suc-eq-zero*:
assumes *integrable*: $\bigwedge i.$ *integrable* *M* (*X i*)
and $\bigwedge i.$ *AE* ξ *in* *M*. *0* = *cond-exp* *M* (*F i*) ($\lambda \xi. X$ (*Suc i*) ξ - *X i* ξ) ξ
shows *nat-martingale* *M F X*
proof (*intro martingale-nat integrable*)
fix *i*
show *AE* ξ *in* *M*. *X i* ξ = *cond-exp* *M* (*F i*) (*X* (*Suc i*)) ξ **using** *cond-exp-diff*[*OF*
integrable(1,1), *of i Suc i i*] *cond-exp-F-meas*[*OF integrable adapted, of i*] *assms*(2)[*of*
i] **by** *fastforce*
qed

0.23 Discrete Time Submartingales

lemma (*in nat-submartingale*) *predictable-ge-zero*:
assumes *nat-predictable-process* *M F X*
shows *AE* ξ *in* *M*. *X i* ξ $\geq X$ *0* ξ
proof (*induction i*)
case *0*
then show ?*case* **by** (*simp add: bot-nat-def*)
next
case (*Suc i*)
interpret *S*: *nat-adapted-process* *M F* $\lambda i.$ *X* (*Suc i*) **by** (*intro nat-predictable-process.adapted-Suc*
assms)
show ?*case* **using** *Suc S.adapted*[*of i*] *submartingale-property*[*OF - le-SucI, of i*]
sigma-finite-subalgebra.cond-exp-F-meas[*OF - integrable, of F i Suc i*] **by** *fastforce*
qed

lemma (*in nat-sigma-finite-adapted-process-order*) *submartingale-of-set-integral-le-Suc*:
assumes *integrable*: $\bigwedge i.$ *integrable* *M* (*X i*)
and $\bigwedge A i.$ *A* $\in F$ *i* \implies *set-lebesgue-integral* *M A* (*X i*) \leq *set-lebesgue-integral*
M A (*X* (*Suc i*))
shows *nat-submartingale* *M F X*
proof (*intro nat-submartingale.intro submartingale-of-set-integral-le*)
fix *i j A* **assume** *asm*: *i* $\leq j$ *A* \in *sets* (*F i*)
show *set-lebesgue-integral* *M A* (*X i*) \leq *set-lebesgue-integral* *M A* (*X j*) **using**


```

asm
proof (induction j - i arbitrary: i j)
  case 0
  then show ?case using asm by simp
next
case (Suc n)
  hence *: n = j - Suc i by linarith
  have Suc i ≤ j using Suc(2,3) by linarith
  thus ?case using sets-F-mono[OF - le-SucI] Suc(4) Suc(1)[OF *] by (auto
intro: assms(2)[THEN order-trans])
qed
qed (simp add: integrable)

lemma (in nat-sigma-finite-adapted-process-order) submartingale-nat:
  assumes integrable:  $\bigwedge i. \text{integrable } M (X i)$ 
  and  $\bigwedge i. AE \xi \text{ in } M. X i \xi \leq \text{cond-exp } M (F i) (X (Suc i)) \xi$ 
  shows nat-submartingale M F X
  using subalg integrable assms(2)
  by (intro submartingale-of-set-integral-le-Suc ord-le-eq-trans[OF set-integral-mono-AE-banach
cond-exp-set-integral[symmetric]], simp)
  (meson in-mono integrable-mult-indicator set-integrable-def subalgebra-def, me-
son integrable-cond-exp in-mono integrable-mult-indicator set-integrable-def subal-
gebra-def, fast+)

lemma (in nat-sigma-finite-adapted-process-order) submartingale-of-cond-exp-diff-Suc-nonneg:
  assumes integrable:  $\bigwedge i. \text{integrable } M (X i)$ 
  and  $\bigwedge i. AE \xi \text{ in } M. 0 \leq \text{cond-exp } M (F i) (\lambda \xi. X (Suc i) \xi - X i \xi) \xi$ 
  shows nat-submartingale M F X
proof (intro submartingale-nat integrable)
  fix i
  show  $AE \xi \text{ in } M. X i \xi \leq \text{cond-exp } M (F i) (X (Suc i)) \xi$  using cond-exp-diff[OF
integrable(1,1), of i Suc i i] cond-exp-F-meas[OF integrable adapted, of i] assms(2)[of
i] by fastforce
qed

lemma (in nat-submartingale) partial-sum-scaleR:
  assumes nat-adapted-process M F C  $\bigwedge i. AE \xi \text{ in } M. 0 \leq C i \xi \bigwedge i. AE \xi \text{ in } M. C i \xi \leq R$ 
  shows nat-submartingale M F ( $\lambda n \xi. \sum i < n. C i \xi *_R (X (Suc i) \xi - X i \xi)$ )
proof -
  interpret C: nat-adapted-process M F C by (rule assms)
  interpret C': nat-adapted-process M F  $\lambda i \xi. C (i - 1) \xi *_R (X i \xi - X (i - 1) \xi)$ 
  by (intro nat-adapted-process.intro adapted-process.scaleR-right adapted-process.diff,
unfold-locales) (auto intro: adaptedD C.adaptedD)+
  interpret C'': nat-adapted-process M F  $\lambda n \xi. \sum i < n. C i \xi *_R (X (Suc i) \xi - X i \xi)$  by (rule C'.partial-sum-Suc[unfolded diff-Suc-1])
  interpret S: nat-sigma-finite-adapted-process-order M F ( $\lambda n \xi. \sum i < n. C i \xi *_R (X (Suc i) \xi - X i \xi)$ ) ..
  have integrable M ( $\lambda x. C i x *_R (X (Suc i) x - X i x)$ ) for i using assms(2,3)[of

```

$i]$ **by** (intro Bochner-Integration.integrable-bound[OF integrable-scaleR-right, OF Bochner-Integration.integrable-diff, OF integrable(1,1), of R Suc i i]) (auto simp add: mult-mono)

moreover have $AE \xi$ in M . $0 \leq \text{cond-exp } M (F i) (\lambda \xi. (\sum_{i < \text{Suc } i} C i \xi *_R (X (\text{Suc } i) \xi - X i \xi)) - (\sum_{i < i} C i \xi *_R (X (\text{Suc } i) \xi - X i \xi))) \xi$ **for** i

using sigma-finite-subalgebra.cond-exp-measurable-scaleR[OF - calculation - C.adapted, of i]

cond-exp-diff-nonneg[OF - le-SucI, OF - order.refl, of i] **assms**(2,3)[of i]

by (fastforce simp add: scaleR-nonneg-nonneg integrable)

ultimately show ?thesis **by** (intro S.submartingale-of-cond-exp-diff-Suc-nonneg Bochner-Integration.integrable-sum, blast+)

qed

lemma (in nat-submartingale) partial-sum-scaleR':

assumes nat-predictable-process $M F C \wedge i$. $AE \xi$ in M . $0 \leq C i \xi \wedge i$. $AE \xi$ in M . $C i \xi \leq R$

shows nat-submartingale $M F (\lambda n \xi. \sum_{i < n} C (\text{Suc } i) \xi *_R (X (\text{Suc } i) \xi - X i \xi))$

proof –

interpret C : nat-predictable-process $M F C$ **by** (rule assms)

interpret $\text{Suc-}C$: nat-adapted-process $M F \lambda i. C (\text{Suc } i)$ **using** C.adapted-Suc .

show ?thesis **by** (intro partial-sum-scaleR[of - R] assms) (intro locales)

qed

0.24 Discrete Time Supermartingales

lemma (in nat-supermartingale) predictable-le-zero:

assumes nat-predictable-process $M F X$

shows $AE \xi$ in M . $X i \xi \leq X 0 \xi$

proof (induction i)

case 0

then show ?case **by** (simp add: bot-nat-def)

next

case (Suc i)

interpret S : nat-adapted-process $M F \lambda i. X (\text{Suc } i)$ **by** (intro nat-predictable-process.adapted-Suc assms)

show ?case **using** Suc S.adapted[of i] supermartingale-property[OF - le-SucI, of i] sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable, of F i Suc i] **by** fastforce

qed

lemma (in nat-sigma-finite-adapted-process-order) supermartingale-of-set-integral-ge-Suc:

assumes integrable: $\bigwedge i$. integrable $M (X i)$

and $\bigwedge A i$. $A \in F i \implies \text{set-lebesgue-integral } M A (X (\text{Suc } i)) \leq \text{set-lebesgue-integral } M A (X i)$

shows nat-supermartingale $M F X$

proof –

interpret -: adapted-process $M F 0 - X$ **by** (rule uminus)

interpret uminus-X: nat-sigma-finite-adapted-process-order $M F - X$..

note * = set-integral-uminus[unfolded set-integrable-def, OF integrable-mult-indicator[OF

```

- integrable]]
  have nat-supermartingale M F (-( - X))
    using ord-eq-le-trans[OF * ord-le-eq-trans[OF le-imp-neg-le[OF assms(2)] * [symmetric]]]
  subalgebra
    by (intro nat-supermartingale.intro submartingale.uminus nat-submartingale.axioms
  uminus-X.submartingale-of-set-integral-le-Suc)
    (clarsimp simp add: fun-Compl-def subalgebra-def integrable | fastforce)+
  thus ?thesis unfolding fun-Compl-def by simp
qed

```

lemma (in nat-sigma-finite-adapted-process-order) supermartingale-nat:

```

  assumes integrable:  $\bigwedge i. \text{integrable } M \ (X \ i)$ 
    and  $\bigwedge i. AE \ \xi \text{ in } M. X \ i \ \xi \geq \text{cond-exp } M \ (F \ i) \ (X \ (Suc \ i)) \ \xi$ 
    shows nat-supermartingale M F X

```

proof -

```

  interpret -: adapted-process M F 0 -X by (rule uminus)
  interpret uminus-X: nat-sigma-finite-adapted-process-order M F -X ..
  have AE  $\xi$  in M.  $- X \ i \ \xi \leq \text{cond-exp } M \ (F \ i) \ (\lambda x. - X \ (Suc \ i) \ x) \ \xi$  for i using
  assms(2) cond-exp-uminus[OF integrable, of i Suc i] by force
  hence nat-supermartingale M F (-( - X)) by (intro nat-supermartingale.intro
  submartingale.uminus nat-submartingale.axioms uminus-X.submartingale-nat) (auto
  simp add: fun-Compl-def integrable)
  thus ?thesis unfolding fun-Compl-def by simp
qed

```

lemma (in nat-sigma-finite-adapted-process-order) supermartingale-of-cond-exp-diff-Suc-nonneg:

```

  assumes integrable:  $\bigwedge i. \text{integrable } M \ (X \ i)$ 
    and  $\bigwedge i. AE \ \xi \text{ in } M. 0 \leq \text{cond-exp } M \ (F \ i) \ (\lambda \xi. X \ i \ \xi - X \ (Suc \ i) \ \xi) \ \xi$ 
    shows nat-supermartingale M F X

```

proof (intro supermartingale-nat integrable)

fix i

```

  show AE  $\xi$  in M.  $X \ i \ \xi \geq \text{cond-exp } M \ (F \ i) \ (X \ (Suc \ i)) \ \xi$  using cond-exp-diff[OF
  integrable(1,1), of i i Suc i] cond-exp-F-meas[OF integrable adapted, of i] assms(2)[of
  i] by fastforce

```

qed

end