

On the Formalization of Martingales

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Abstract

This thesis presents a formalization of martingales in arbitrary Banach spaces using Isabelle/HOL. We begin by examining formalizations in prominent proof repositories and extend the definition of the conditional expectation operator from real numbers to general Banach spaces, drawing inspiration from prior work. We define filtered measure spaces, adapted, progressively measurable, and predictable processes and rigorously formalize martingales, submartingales, and supermartingales. Additionally, our contributions expand the scope of Bochner integration techniques to general Banach spaces and introduce additional lemmas and induction schemes for integrable functions. Our formalization provides a robust framework for future developments within the theory of stochastic processes. Furthermore, we provide an example demonstrating when a coin toss constitutes a martingale, submartingale, or a supermartingale.

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```

theory Measure-Space-Supplement
  imports HOL-Analysis.Measure-Space
begin

```

1 Supplementary Lemmas for Measure Spaces

1.1 Sigma Algebra Generated by a Family of Functions

definition *family-vimage-algebra* :: 'a set \Rightarrow ('a \Rightarrow 'b) set \Rightarrow 'b measure \Rightarrow 'a measure **where**
family-vimage-algebra Ω S $M \equiv \text{sigma } \Omega (\bigcup f \in S. \{f - 'A \cap \Omega \mid A. A \in M\})$

lemma *family-vimage-algebra-singleton*: *family-vimage-algebra* Ω $\{f\}$ $M = \text{vimage-algebra } \Omega f M$ **unfolding** *family-vimage-algebra-def* *vimage-algebra-def* **by** *simp*

lemma

shows *sets-family-vimage-algebra*: *sets* (*family-vimage-algebra* Ω S M) = *sigma-sets* $\Omega (\bigcup f \in S. \{f - 'A \cap \Omega \mid A. A \in M\})$

and *space-family-vimage-algebra[simp]*: *space* (*family-vimage-algebra* Ω S M) = Ω

by (*auto simp add: family-vimage-algebra-def sets-measure-of-conv space-measure-of-conv*)

lemma *measurable-family-vimage-algebra*:

assumes $f \in S$ $f \in \Omega \rightarrow \text{space } M$

shows $f \in \text{family-vimage-algebra } \Omega S M \rightarrow_M M$

using *assms* **by** (*intro measurableI, auto simp add: sets-family-vimage-algebra*)

lemma *measurable-family-vimage-algebra-singleton*:

assumes $f \in \Omega \rightarrow \text{space } M$

shows $f \in \text{family-vimage-algebra } \Omega \{f\} M \rightarrow_M M$

using *assms* *measurable-family-vimage-algebra* **by** *blast*

lemma *measurable-family-iff-sets*:

shows $(S \subseteq N \rightarrow_M M) \longleftrightarrow S \subseteq \text{space } N \rightarrow \text{space } M \wedge \text{family-vimage-algebra } (\text{space } N) S M \subseteq N$

proof (*standard, goal-cases*)

case 1

hence *subset*: $S \subseteq \text{space } N \rightarrow \text{space } M$ **using** *measurable-space* **by** *fast*

have $\{f - 'A \cap \text{space } N \mid A. A \in M\} \subseteq N$ **if** $f \in S$ **for** f **using** *measurable-iff-sets[unfolded family-vimage-algebra-singleton[symmetric], of f]* 1 *subset that*
by (*fastforce simp add: sets-family-vimage-algebra*)

then show ?*case* **unfolding** *sets-family-vimage-algebra* **using** *sets.sigma-algebra-axioms*
by (*simp add: subset, intro sigma-algebra.sigma-sets-subset, blast+*)

next

case 2

hence *subset*: $S \subseteq \text{space } N \rightarrow \text{space } M$ **by** *simp*

show ?*case*

```

proof (standard, goal-cases)
  case (1 x)
    have family-vimage-algebra (space N) {x} M  $\subseteq$  N by (metis (no-types, lifting)
1 2 sets-family-vimage-algebra SUP-le-iff sigma-sets-le-sets-iff singletonD)
    thus ?case using measurable-iff-sets[unfolded family-vimage-algebra-singleton[symmetric]]
subset[THEN subsetD, OF 1] by fast
  qed
qed

lemma family-vimage-algebra-diff:
  shows family-vimage-algebra  $\Omega$  S M = sigma  $\Omega$  (sets (family-vimage-algebra  $\Omega$ 
(S - I) M)  $\cup$  family-vimage-algebra  $\Omega$  (S  $\cap$  I) M)
  using sets.space-closed space-measure-of-conv
  unfolding family-vimage-algebra-def sets-family-vimage-algebra
  by (intro sigma-eqI, blast, fastforce)
    (intro sigma-sets-eqI, blast, simp add: sets-measure-of-conv split: if-splits,
meson Diff-subset Sup-subset-mono in-mono inf-sup-ord(1) sigma-sets-subseteq
subset-image-iff, fastforce+)

end
theory Elementary-Metric-Spaces-Supplement
  imports HOL-Analysis.Elementary-Metric-Spaces
begin

```

2 Supplementary Lemmas for Elementary Metric Spaces

2.1 Diameter Lemma

```

lemma diameter-comp-strict-mono:
  fixes s :: nat  $\Rightarrow$  'a :: metric-space
  assumes strict-mono r bounded {s i | i. r n  $\leq$  i}
  shows diameter {s (r i) | i. n  $\leq$  i}  $\leq$  diameter {s i | i. r n  $\leq$  i}
proof (rule diameter-subset)
  show {s (r i) | i. n  $\leq$  i}  $\subseteq$  {s i | i. r n  $\leq$  i} using assms(1) monotoneD
strict-mono-mono by fastforce
qed (intro assms(2))

lemma diameter-bounded-bound':
  fixes S :: 'a :: metric-space set
  assumes S: bdd-above (case-prod dist ' (S $\times$ S)) x  $\in$  S y  $\in$  S
  shows dist x y  $\leq$  diameter S
proof -
  have (x,y)  $\in$  S $\times$ S using S by auto
  then have dist x y  $\leq$  (SUP (x,y) $\in$ S $\times$ S. dist x y) by (rule cSUP-upper2[OF
assms(1)]) simp
  with  $\langle x \in S \rangle$  show ?thesis by (auto simp: diameter-def)
qed

```

```

lemma bounded-imp-dist-bounded:
  assumes bounded (range s)
  shows bounded (( $\lambda(i, j). \text{dist } (s \ i) \ (s \ j)$ ) ‘ ( $\{n..\} \times \{n..\}$ ))
  using bounded-dist-comp[OF bounded-fst bounded-snd, OF bounded-Times(1,1)[OF
assms(1,1)]] by (rule bounded-subset, force)

lemma cauchy-iff-diameter-tends-to-zero-and-bounded:
  fixes s :: nat  $\Rightarrow$  'a :: metric-space
  shows Cauchy s  $\longleftrightarrow$  (( $\lambda n. \text{diameter } \{s \ i \mid i. i \geq n\}$ )  $\longrightarrow$  0  $\wedge$  bounded (range
s))
proof –
  have  $\{s \ i \mid i. N \leq i\} \neq \{\}$  for N by blast
  hence diameter-SUP: diameter  $\{s \ i \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s \ i) \ (s \ j))$  for N unfolding diameter-def by (auto intro!: arg-cong[of - - Sup])
  show ?thesis
  proof (intro iffI)
    assume asm: Cauchy s
    have  $\exists N. \forall n \geq N. \text{norm } (\text{diameter } \{s \ i \mid i. n \leq i\}) < e$  if e-pos: e > 0 for e
    proof –
      obtain N where dist-less: dist (s n) (s m) < (1/2) * e if n  $\geq$  N m  $\geq$  N
for n m using asm e-pos by (metis Cauchy-def mult-pos-pos zero-less-divide-iff
zero-less-numeral zero-less-one)
      {
        fix r assume r  $\geq$  N
        hence dist (s n) (s m) < (1/2) * e if n  $\geq$  r m  $\geq$  r for n m using dist-less
        that by simp
        hence ( $\text{SUP } (i, j) \in \{r..\} \times \{r..\}. \text{dist } (s \ i) \ (s \ j)$ )  $\leq$  (1/2) * e by (intro
cSup-least) fastforce+
        also have ... < e using e-pos by simp
        finally have diameter  $\{s \ i \mid i. r \leq i\} < e$  using diameter-SUP by presburger
      }
      moreover have diameter  $\{s \ i \mid i. r \leq i\} \geq 0$  for r unfolding diameter-SUP
using bounded-imp-dist-bounded[OF cauchy-imp-bounded, THEN bounded-imp-bdd-above,
OF asm] by (intro cSup-upper2, auto)
      ultimately show ?thesis by auto
    qed
    thus ( $\lambda n. \text{diameter } \{s \ i \mid i. n \leq i\}$ )  $\longrightarrow$  0  $\wedge$  bounded (range s) using
cauchy-imp-bounded[OF asm] by (simp add: LIMSEQ-iff)
  next
    assume asm: ( $\lambda n. \text{diameter } \{s \ i \mid i. n \leq i\}$ )  $\longrightarrow$  0  $\wedge$  bounded (range s)
    have  $\exists N. \forall n \geq N. \forall m \geq N. \text{dist } (s \ n) \ (s \ m) < e$  if e-pos: e > 0 for e
    proof –
      obtain N where diam-less: diameter  $\{s \ i \mid i. r \leq i\} < e$  if r  $\geq$  N for r
using LIMSEQ-D asm(1) e-pos by fastforce
      {
        fix n m assume n  $\geq$  N m  $\geq$  N
        hence dist (s n) (s m) < e using cSUP-lessD[OF bounded-imp-dist-bounded[THEN
bounded-imp-bdd-above], OF - diam-less[unfolded diameter-SUP]] asm by auto
      }

```

```

    thus ?thesis by blast
  qed
  then show Cauchy s by (simp add: Cauchy-def)
  qed
qed
end

```

```

theory Bochner-Integration-Supplement
  imports HOL-Analysis.Bochner-Integration HOL-Analysis.Set-Integral Elementary-Metric-Spaces-Supplement
begin

```

3 Supplementary Lemmas for Bochner Integration

3.1 Integrable Simple Functions

```

lemma integrable-implies-simple-function-sequence:
  fixes f :: 'a  $\Rightarrow$  'b :: {banach, second-countable-topology}
  assumes integrable M f
  obtains s where  $\bigwedge i. \text{simple-function } M (s\ i)$ 
    and  $\bigwedge i. \text{emeasure } M \{y \in \text{space } M. s\ i\ y \neq 0\} \neq \infty$ 
    and  $\bigwedge x. x \in \text{space } M \implies (\lambda i. s\ i\ x) \longrightarrow f\ x$ 
    and  $\bigwedge x\ i. x \in \text{space } M \implies \text{norm } (s\ i\ x) \leq 2 * \text{norm } (f\ x)$ 
proof -
  have f:  $f \in \text{borel-measurable } M (\int^+ x. \text{norm } (f\ x) \partial M) < \infty$  using assms
  unfolding integrable-iff-bounded by auto
  obtain s where s:  $\bigwedge i. \text{simple-function } M (s\ i) \bigwedge x. x \in \text{space } M \implies (\lambda i. s\ i\ x) \longrightarrow f\ x$ 
     $\bigwedge i\ x. x \in \text{space } M \implies \text{norm } (s\ i\ x) \leq 2 * \text{norm } (f\ x)$  using
    borel-measurable-implies-sequence-metric[OF f(1)] unfolding norm-conv-dist by
    metis
  {
    fix i
    have  $(\int^+ x. \text{norm } (s\ i\ x) \partial M) \leq (\int^+ x. \text{ennreal } (2 * \text{norm } (f\ x)) \partial M)$  using
    s by (intro nn-integral-mono, auto)
    also have  $\dots < \infty$  using f by (simp add: nn-integral-cmult ennreal-mult-less-top
    ennreal-mult)
    finally have sbi: Bochner-Integration.simple-bochner-integrable M (s i) using
    s by (intro simple-bochner-integrableI-bounded) auto
    hence  $\text{emeasure } M \{y \in \text{space } M. s\ i\ y \neq 0\} \neq \infty$  by (auto intro: integrableI-simple-bochner-integrable
    simple-bochner-integrable.cases)
  }
  thus ?thesis using that s by blast
qed

```

```

lemma simple-function-indicator-representation:
  fixes f :: 'a  $\Rightarrow$  'b :: {second-countable-topology, banach}
  assumes f: simple-function M f and x:  $x \in \text{space } M$ 

```

shows $f x = (\sum y \in f \text{ ' space } M. \text{ indicator } (f - \text{' } \{y\} \cap \text{ space } M) x *_R y)$
(is $?l = ?r)$
proof –
have $?r = (\sum y \in f \text{ ' space } M.$
 $(\text{if } y = f x \text{ then indicator } (f - \text{' } \{y\} \cap \text{ space } M) x *_R y \text{ else } 0))$ **by** $(\text{auto intro!}:$
 $\text{sum.cong})$
also have $\dots = \text{indicator } (f - \text{' } \{f x\} \cap \text{ space } M) x *_R f x$ **using** assms **by** $(\text{auto}$
 $\text{dest: simple-functionD})$
also have $\dots = f x$ **using** x **by** $(\text{auto simp: indicator-def})$
finally show $?thesis$ **by** auto
qed

lemma *simple-function-indicator-representation-AE*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes f : *simple-function* $M f$
shows $\text{AE } x \text{ in } M. f x = (\sum y \in f \text{ ' space } M. \text{ indicator } (f - \text{' } \{y\} \cap \text{ space } M) x$
 $*_R y)$
by $(\text{metis (mono-tags, lifting) AE-I2 simple-function-indicator-representation } f)$

lemmas *simple-function-scaleR[intro] = simple-function-compose2[where h=(*_R)]*
lemmas *integrable-simple-function = simple-bochner-integrable.intros[THEN has-bochner-integral-simple-bochner]*
 $\text{THEN integrable.intros}$

Induction rule for simple integrable functions.

lemma *integrable-simple-function-induct[consumes 2, case-names cong indicator*
 $\text{add, induct set: simple-function}]$:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes f : *simple-function* $M f$ $\text{emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty$
assumes cong : $\bigwedge f g. \text{ simple-function } M f \implies \text{emeasure } M \{y \in \text{space } M. f y \neq$
 $0\} \neq \infty$
 $\implies \text{simple-function } M g \implies \text{emeasure } M \{y \in \text{space } M. g y \neq$
 $0\} \neq \infty$
 $\implies (\bigwedge x. x \in \text{space } M \implies f x = g x) \implies P f \implies P g$
assumes indicator : $\bigwedge A y. A \in \text{sets } M \implies \text{emeasure } M A < \infty \implies P (\lambda x.$
 $\text{indicator } A x *_R y)$
assumes add : $\bigwedge f g. \text{ simple-function } M f \implies \text{emeasure } M \{y \in \text{space } M. f y \neq$
 $0\} \neq \infty \implies$
 $\text{simple-function } M g \implies \text{emeasure } M \{y \in \text{space } M. g y \neq 0\} \neq$
 $\infty \implies$
 $(\bigwedge z. z \in \text{space } M \implies \text{norm } (f z + g z) = \text{norm } (f z) + \text{norm}$
 $(g z)) \implies$
 $P f \implies P g \implies P (\lambda x. f x + g x)$
shows $P f$
proof –
let $?f = \lambda x. (\sum y \in f \text{ ' space } M. \text{ indicat-real } (f - \text{' } \{y\} \cap \text{ space } M) x *_R y)$
have $f\text{-ae-eq}$: $f x = ?f x$ **if** $x \in \text{space } M$ **for** x **using** *simple-function-indicator-representation[OF*
 $f(1) \text{ that}]$.
moreover have $\text{emeasure } M \{y \in \text{space } M. ?f y \neq 0\} \neq \infty$ **by** $(\text{metis (no-types,}$
 $\text{lifting) Collect-cong calculation } f(2))$

moreover have $P (\lambda x. \sum_{y \in S}. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y)$
 $\text{simple-function } M (\lambda x. \sum_{y \in S}. \text{indicat-real } (f - \{y\} \cap \text{space } M) x$
 $*_R y)$
 $\text{emeasure } M \{y \in \text{space } M. (\sum_{x \in S}. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \neq \infty$
if $S \subseteq f - \{ \text{space } M \}$ **for** S **using** $\text{simple-functionD}(1)[\text{OF } \text{assms}(1),$
 $\text{THEN } \text{rev-finite-subset, OF that}] \text{ that}$
proof (*induction rule: finite-induct*)
case empty
 $\{$
case 1
then show $?case \text{ using indicator}[of \{ \}] \text{ by force}$
next
case 2
then show $?case \text{ by force}$
next
case 3
then show $?case \text{ by force}$
 $\}$
next
case (*insert x F*)
have $(f - \{x\} \cap \text{space } M) \subseteq \{y \in \text{space } M. f y \neq 0\}$ **if** $x \neq 0$ **using that by**
 blast
moreover have $\{y \in \text{space } M. f y \neq 0\} = \text{space } M - (f - \{0\} \cap \text{space } M)$
by blast
moreover have $\text{space } M - (f - \{0\} \cap \text{space } M) \in \text{sets } M$ **using** $\text{simple-functionD}(2)[\text{OF } f(1)] \text{ by blast}$
ultimately have $\text{fin-0: emeasure } M (f - \{x\} \cap \text{space } M) < \infty$ **if** $x \neq 0$
using that by (*metis emeasure-mono f(2) infinity-ennreal-def top.not-eq-extremum top-unique*)
hence $\text{fin-1: emeasure } M \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R$
 $x \neq 0\} \neq \infty$ **if** $x \neq 0$ **by** (*metis (mono-tags, lifting) emeasure-mono f(1) indica-*
 $\text{tor-simps}(2) \text{ linorder-not-less mem-Collect-eq scaleR-eq-0-iff simple-functionD}(2)$
 $\text{subsetI that})$

have $*$: $(\sum_{y \in \text{insert } x F}. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y) = (\sum_{y \in F}. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y) + \text{indicat-real } (f - \{x\} \cap \text{space } M)$
 $x *_R x$ **for** x **by** (*metis (no-types, lifting) Diff-empty Diff-insert0 add commute insert.hyps(1) insert.hyps(2) sum.insert-remove*)
have $**$: $\{y \in \text{space } M. (\sum_{x \in \text{insert } x F}. \text{indicat-real } (f - \{x\} \cap \text{space } M) y$
 $*_R x) \neq 0\} \subseteq \{y \in \text{space } M. (\sum_{x \in F}. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x)$
 $\neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\}$ **unfolding**
 $*$ **by** fastforce
 $\{$
case 1
hence $x: x \in f - \{ \text{space } M \}$ **and** $F: F \subseteq f - \{ \text{space } M \}$ **by auto**
show $?case$
proof (*cases x = 0*)
case True


```

    then show ?thesis unfolding * using insert(3)[OF F] by simp
next
  case False
  have norm-argument: norm (( $\sum_{y \in F} \text{indicat-real } (f - \{y\} \cap \text{space } M) z$ 
 $*_R y) + \text{indicat-real } (f - \{x\} \cap \text{space } M) z *_R x) = \text{norm } (\sum_{y \in F} \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) + \text{norm } (\text{indicat-real } (f - \{x\} \cap \text{space } M) z *_R x)$ 
  if  $z: z \in \text{space } M$  for  $z$ 
  proof (cases  $f z = x$ )
  case True
  have indicat-real  $(f - \{y\} \cap \text{space } M) z *_R y = 0$  if  $y \in F$  for  $y$  using
  True insert(2)  $z$  that 1 unfolding indicator-def by force
  hence  $(\sum_{y \in F} \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) = 0$  by (meson
  sum.neutral)
  then show ?thesis by force
  next
  case False
  then show ?thesis by force
qed
show ?thesis using False simple-functionD(2)[OF  $f(1)$ ] insert(3,5)[OF F]
simple-function-scaleR fin-0 fin-1 by (subst *, subst add, subst simple-function-sum)
(blast intro: norm-argument indicator)+
qed
next
  case 2
  hence  $x: x \in f \text{ ' space } M$  and  $F: F \subseteq f \text{ ' space } M$  by auto
  show ?case
  proof (cases  $x = 0$ )
  case True
  then show ?thesis unfolding * using insert(4)[OF F] by simp
  next
  case False
  then show ?thesis unfolding * using insert(4)[OF F] simple-functionD(2)[OF
 $f(1)$ ] by fast
  qed
next
  case 3
  hence  $x: x \in f \text{ ' space } M$  and  $F: F \subseteq f \text{ ' space } M$  by auto
  show ?case
  proof (cases  $x = 0$ )
  case True
  then show ?thesis unfolding * using insert(5)[OF F] by simp
  next
  case False
  have emeasure  $M \{y \in \text{space } M. (\sum_{x \in \text{insert } x F} \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \leq \text{emeasure } M (\{y \in \text{space } M. (\sum_{x \in F} \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\})$ 
  using ** simple-functionD(2)[OF insert(4)[OF F]] simple-functionD(2)[OF
 $f(1)$ ] by (intro emeasure-mono, force+)

```

also have $\dots \leq \text{emeasure } M \{y \in \text{space } M. (\sum x \in F. \text{indicat-real } (f - \cdot \{x\} \cap \text{space } M) y *_R x) \neq 0\} + \text{emeasure } M \{y \in \text{space } M. \text{indicat-real } (f - \cdot \{x\} \cap \text{space } M) y *_R x \neq 0\}$
using $\text{simple-functionD}(2)[OF \text{ insert}(4)[OF F]] \text{ simple-functionD}(2)[OF f(1)]$ **by** $(\text{intro } \text{emeasure-subadditive}, \text{force+})$
also have $\dots < \infty$ **using** $\text{insert}(5)[OF F] \text{ fin-1}[OF \text{ False}]$ **by** $(\text{simp add: less-top})$
finally show $?thesis$ **by** simp
qed
}
qed
moreover have $\text{simple-function } M (\lambda x. \sum y \in f \cdot \text{space } M. \text{indicat-real } (f - \cdot \{y\} \cap \text{space } M) x *_R y)$ **using** calculation **by** blast
moreover have $P (\lambda x. \sum y \in f \cdot \text{space } M. \text{indicat-real } (f - \cdot \{y\} \cap \text{space } M) x *_R y)$ **using** calculation **by** blast
ultimately show $?thesis$ **by** $(\text{intro } \text{cong}[OF - \cdot f(1,2)], \text{blast}, \text{presburger+})$
qed

Induction rule for non-negative simple integrable functions

lemma $\text{integrable-simple-function-induct-nn}[\text{consumes } 3, \text{case-names } \text{cong } \text{indicator } \text{add}, \text{induct set: } \text{simple-function}]$:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}, \text{linorder-topology}, \text{ordered-real-vector}\}$

assumes f : $\text{simple-function } M f$ $\text{emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty \wedge x \in \text{space } M \longrightarrow f x \geq 0$

assumes cong : $\bigwedge f g. \text{simple-function } M f \implies \text{emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty \implies (\bigwedge x. x \in \text{space } M \implies f x \geq 0) \implies \text{simple-function } M g \implies \text{emeasure } M \{y \in \text{space } M. g y \neq 0\} \neq \infty \implies (\bigwedge x. x \in \text{space } M \implies g x \geq 0) \implies (\bigwedge x. x \in \text{space } M \implies f x = g x) \implies P f \implies P g$

assumes indicator : $\bigwedge A y. y \geq 0 \implies A \in \text{sets } M \implies \text{emeasure } M A < \infty \implies P (\lambda x. \text{indicator } A x *_R y)$

assumes add : $\bigwedge f g. (\bigwedge x. x \in \text{space } M \implies f x \geq 0) \implies \text{simple-function } M f \implies \text{emeasure } M \{y \in \text{space } M. f y \neq 0\} \neq \infty \implies$

$(\bigwedge x. x \in \text{space } M \implies g x \geq 0) \implies \text{simple-function } M g \implies \text{emeasure } M \{y \in \text{space } M. g y \neq 0\} \neq \infty \implies$
 $(\bigwedge z. z \in \text{space } M \implies \text{norm } (f z + g z) = \text{norm } (f z) + \text{norm } (g z)) \implies$

$P f \implies P g \implies P (\lambda x. f x + g x)$

shows $P f$

proof–

let $?f = \lambda x. (\sum y \in f \cdot \text{space } M. \text{indicat-real } (f - \cdot \{y\} \cap \text{space } M) x *_R y)$

have $f\text{-ae-eq}$: $f x = ?f x$ **if** $x \in \text{space } M$ **for** x **using** $\text{simple-function-indicator-representation}[OF f(1) \text{ that}]$.

moreover have $\text{emeasure } M \{y \in \text{space } M. ?f y \neq 0\} \neq \infty$ **by** $(\text{metis } (\text{no-types}, \text{lifting}) \text{ Collect-cong } \text{calculation } f(2))$

moreover have $P (\lambda x. \sum y \in S. \text{indicat-real } (f - \cdot \{y\} \cap \text{space } M) x *_R y)$

$\text{simple-function } M (\lambda x. \sum y \in S. \text{indicat-real } (f - \cdot \{y\} \cap \text{space } M) x$

$*_R y)$

$\text{emeasure } M \{y \in \text{space } M. (\sum x \in S. \text{indicat-real } (f - \cdot \{x\} \cap \text{space$

$M) y *_R x) \neq 0\} \neq \infty$
 $\bigwedge x. x \in \text{space } M \implies 0 \leq (\sum y \in S. \text{indicat-real } (f - \{y\} \cap \text{space } M)$
 $x *_R y)$
if $S \subseteq f - \{y\} \cap \text{space } M$ **for** S **using** $\text{simple-functionD}(1)[\text{OF } \text{assms}(1),$
THEN $\text{rev-finite-subset}, \text{OF that}]$ **that**
proof ($\text{induction rule: finite-subset-induct}$)
case empty
{
case 1
then show $?case$ **using** $\text{indicator}[of 0 \{y\}]$ **by force**
next
case 2
then show $?case$ **by force**
next
case 3
then show $?case$ **by force**
next
case 4
then show $?case$ **by force**
}
next
case ($\text{insert } x F$)
have $(f - \{x\} \cap \text{space } M) \subseteq \{y \in \text{space } M. f y \neq 0\}$ **if** $x \neq 0$ **using that by**
 blast
moreover have $\{y \in \text{space } M. f y \neq 0\} = \text{space } M - (f - \{0\} \cap \text{space } M)$
by blast
moreover have $\text{space } M - (f - \{0\} \cap \text{space } M) \in \text{sets } M$ **using** $\text{simple-functionD}(2)[\text{OF } f(1)]$ **by blast**
ultimately have $\text{fin-0: emeasure } M (f - \{x\} \cap \text{space } M) < \infty$ **if** $x \neq 0$
using that by ($\text{metis emeasure-mono } f(2) \text{ infinity-ennreal-def top.not-eq-extremum}$
 top-unique)
**hence fin-1: emeasure } M \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R
 $x \neq 0\} \neq \infty$ **if** $x \neq 0$ **by** ($\text{metis (mono-tags, lifting) emeasure-mono } f(1) \text{ indica-}$
 $\text{tor-simps}(2) \text{ linorder-not-less mem-Collect-eq scaleR-eq-0-iff simple-functionD}(2)$
 subsetI that)

have nonneg-x: } x \geq 0 **using insert } f **by blast
**have *: } (\sum y \in \text{insert } x F. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y) =
 $(\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) x *_R y) + \text{indicat-real } (f - \{x\} \cap$
 $\text{space } M) x *_R x$ **for** x **by** ($\text{metis (no-types, lifting) add.commute insert.hyps}(1)$
 $\text{insert.hyps}(4) \text{ sum.insert}$)
**have **: } \{y \in \text{space } M. (\sum x \in \text{insert } x F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y
 $*_R x) \neq 0\} \subseteq \{y \in \text{space } M. (\sum x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x)$
 $\neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\}$ **unfolding**
 $*$ **by fastforce**
{
case 1
show $?case$
proof ($\text{cases } x = 0$)**********

```

    case True
    then show ?thesis unfolding * using insert by simp
next
    case False
    have norm-argument: norm (( $\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) z$ 
 $*_R y) + \text{indicat-real } (f - \{x\} \cap \text{space } M) z *_R x) = \text{norm } (\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) + \text{norm } (\text{indicat-real } (f - \{x\} \cap \text{space } M) z *_R x)$ 
    if  $z: z \in \text{space } M$  for  $z$ 
    proof (cases  $f z = x$ )
    case True
    have  $\text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y = 0$  if  $y \in F$  for  $y$  using
    True insert  $z$  that 1 unfolding indicator-def by force
    hence  $(\sum y \in F. \text{indicat-real } (f - \{y\} \cap \text{space } M) z *_R y) = 0$  by (meson
    sum.neutral)
    thus ?thesis by force
    qed (force)
    show ?thesis using False fin-0 fin-1 f norm-argument by (subst *, subst add,
    presburger add: insert, intro insert, intro insert, fastforce simp add: indicator-def
    intro!: insert(2) f(3), auto intro!: indicator insert nonneg-x)
    qed
next
    case 2
    show ?case
    proof (cases  $x = 0$ )
    case True
    then show ?thesis unfolding * using insert by simp
    next
    case False
    then show ?thesis unfolding * using insert simple-functionD(2)[OF f(1)]
by fast
    qed
next
    case 3
    show ?case
    proof (cases  $x = 0$ )
    case True
    then show ?thesis unfolding * using insert by simp
    next
    case False
    have  $\text{emeasure } M \{y \in \text{space } M. (\sum x \in \text{insert } x F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \leq \text{emeasure } M \{y \in \text{space } M. (\sum x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} \cup \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\}$ 
    using ** simple-functionD(2)[OF insert(6)] simple-functionD(2)[OF f(1)]
    insert.IH(2) by (intro emeasure-mono, blast, simp)
    also have  $\dots \leq \text{emeasure } M \{y \in \text{space } M. (\sum x \in F. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x) \neq 0\} + \text{emeasure } M \{y \in \text{space } M. \text{indicat-real } (f - \{x\} \cap \text{space } M) y *_R x \neq 0\}$ 
    using simple-functionD(2)[OF insert(6)] simple-functionD(2)[OF f(1)]

```

by (*intro emeasure-subadditive, force+*)
 also have $\dots < \infty$ **using** *insert*(γ) *fin-1*[*OF False*] **by** (*simp add: less-top*)
 finally show *?thesis* **by** *simp*
 qed
next
 case (ξ)
 thus *?case using insert nonneg-x f(3)* **by** (*auto simp add: scaleR-nonneg-nonneg*
intro: sum-nonneg)
 qed
moreover have *simple-function M* ($\lambda x. \sum y \in f^{-1} \text{space } M. \text{indicat-real } (f^{-1} \{y\} \cap \text{space } M) x *_R y$) **using** *calculation by blast*
moreover have *P* ($\lambda x. \sum y \in f^{-1} \text{space } M. \text{indicat-real } (f^{-1} \{y\} \cap \text{space } M) x *_R y$) **using** *calculation by blast*
moreover have $\bigwedge x. x \in \text{space } M \implies 0 \leq f x$ **using** *f(3) by simp*
ultimately show *?thesis by (intro cong[OF - - - f(1,2)]), blast, blast, fast*
presburger+
qed

lemma *finite-nn-integral-imp-ae-finite:*

fixes $f :: 'a \Rightarrow \text{ennreal}$
assumes $f \in \text{borel-measurable } M$ ($\int^+ x. f x \, \partial M < \infty$)
shows *AE x in M. f x < ∞*
proof (*rule ccontr, goal-cases*)
case 1
let $?A = \text{space } M \cap \{x. f x = \infty\}$
have $*$: *emeasure M ?A > 0 using 1 assms(1) by (metis (mono-tags, lifting) assms(2) eventually-mono infinity-ennreal-def nn-integral-not-eq-infinite top.not-eq-extremum)*
have ($\int^+ x. f x * \text{indicator } ?A x \, \partial M$) = ($\int^+ x. \infty * \text{indicator } ?A x \, \partial M$) **by**
(metis (mono-tags, lifting) indicator-inter-arith indicator-simps(2) mem-Collect-eq mult-eq-0-iff)
also have $\dots = \infty * \text{emeasure } M ?A$ **using** *assms(1) by (intro nn-integral-cmult-indicator, simp)*
also have $\dots = \infty$ **using** $*$ **by** *fastforce*
finally have ($\int^+ x. f x * \text{indicator } ?A x \, \partial M$) = ∞ .
moreover have ($\int^+ x. f x * \text{indicator } ?A x \, \partial M$) \leq ($\int^+ x. f x \, \partial M$) **by** (*intro nn-integral-mono, simp add: indicator-def*)
ultimately show *?case using assms(2) by simp*
qed

Convergence in L1-Norm implies existence of a subsequence which converges almost everywhere. This lemma is easier to use than the existing one in *HOL-Analysis.Bochner-Integration*

lemma *cauchy-L1-AE-cauchy-subseq:*

fixes $s :: \text{nat} \Rightarrow 'a \Rightarrow 'b :: \{\text{banach, second-countable-topology}\}$
assumes [*measurable*]: $\bigwedge n. \text{integrable } M (s \, n)$
and $\bigwedge e. e > 0 \implies \exists N. \forall i \geq N. \forall j \geq N. \text{LINT } x | M. \text{norm } (s \, i \, x - s \, j \, x) < e$
obtains r **where** *strict-mono r AE x in M. Cauchy* ($\lambda i. s \, (r \, i) \, x$)
proof–

have $\exists r. \forall n. (\forall i \geq r n. \forall j \geq r n. \text{LINT } x | M. \text{norm } (s \ i \ x - s \ j \ x) < (1 / 2) \wedge n) \wedge (r \ (\text{Suc } n) > r \ n)$
proof (intro dependent-nat-choice, goal-cases)
case 1
then show ?case **using** assms(2) **by** presburger
next
case (2 x n)
obtain N **where** *: $\text{LINT } x | M. \text{norm } (s \ i \ x - s \ j \ x) < (1 / 2) \wedge \text{Suc } n$ **if** $i \geq N \ j \geq N$ **for** i j **using** assms(2)[of (1 / 2) \wedge Suc n] **by** auto
{
fix i j **assume** $i \geq \max N \ (\text{Suc } x) \ j \geq \max N \ (\text{Suc } x)$
hence $\text{integral}^L M \ (\lambda x. \text{norm } (s \ i \ x - s \ j \ x)) < (1 / 2) \wedge \text{Suc } n$ **using** * **by** fastforce
}
then show ?case **by** fastforce
qed
then obtain r **where** strict-mono: strict-mono r **and** $\forall i \geq r n. \forall j \geq r n. \text{LINT } x | M. \text{norm } (s \ i \ x - s \ j \ x) < (1 / 2) \wedge n$ **for** n **using** strict-mono-Suc-iff **by** blast
hence r-is: $\text{LINT } x | M. \text{norm } (s \ (r \ (\text{Suc } n)) \ x - s \ (r \ n) \ x) < (1 / 2) \wedge n$ **for** n **by** (simp add: strict-mono-leD)

define g **where** $g = (\lambda n \ x. (\sum i \leq n. \text{ennreal } (\text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x))))$
define g' **where** $g' = (\lambda x. \sum i. \text{ennreal } (\text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x)))$

have integrable-g: $(\int^+ x. g \ n \ x \ \partial M) < 2$ **for** n
proof -
have $(\int^+ x. g \ n \ x \ \partial M) = (\int^+ x. (\sum i \leq n. \text{ennreal } (\text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x))) \ \partial M)$ **using** g-def **by** simp
also have ... = $(\sum i \leq n. (\int^+ x. \text{ennreal } (\text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x)) \ \partial M))$ **by** (intro nn-integral-sum, simp)
also have ... = $(\sum i \leq n. \text{LINT } x | M. \text{norm } (s \ (r \ (\text{Suc } i)) \ x - s \ (r \ i) \ x))$
unfolding dist-norm **using** assms(1) **by** (subst nn-integral-eq-integral[OF integrable-norm], auto)
also have ... < $\text{ennreal } (\sum i \leq n. (1 / 2) \wedge i)$ **by** (intro ennreal-lessI[OF sum-pos sum-strict-mono[OF finite-atMost - r-is]], auto)
also have ... $\leq \text{ennreal } 2$ **unfolding** sum-gp0[$\text{of } 1 / 2 \ n$] **by** (intro ennreal-leI, auto)
finally show $(\int^+ x. g \ n \ x \ \partial M) < 2$ **by** simp
qed

have integrable-g': $(\int^+ x. g' \ x \ \partial M) \leq 2$
proof -
have incseq $(\lambda n. g \ n \ x)$ **for** x **by** (intro incseq-SucI, auto simp add: g-def ennreal-leI)
hence convergent $(\lambda n. g \ n \ x)$ **for** x **unfolding** convergent-def **using** LIMSEQ-incseq-SUP **by** fast
hence $(\lambda n. g \ n \ x) \longrightarrow g' \ x$ **for** x **unfolding** g-def g'-def **by** (intro summable-iff-convergent'[THEN iffD2, THEN summable-LIMSEQ], blast)

hence $(\int^+ x. g' x \partial M) = (\int^+ x. \liminf (\lambda n. g n x) \partial M)$ **by** (*metis lim-imp-Liminf trivial-limit-sequentially*)
also have $\dots \leq \liminf (\lambda n. \int^+ x. g n x \partial M)$ **by** (*intro nn-integral-liminf, simp add: g-def*)
also have $\dots \leq \liminf (\lambda n. 2)$ **using** *integrable-g* **by** (*intro Liminf-mono*) (*simp add: order-le-less*)
also have $\dots = 2$ **using** *sequentially-bot tendsto-iff-Liminf-eq-Limsup* **by** *blast*
finally show *?thesis* .
qed
hence *AE x in M. g' x < ∞* **by** (*intro finite-nn-integral-imp-ae-finite*) (*auto simp add: order-le-less-trans g'-def*)
moreover have *summable (λi. norm (s (r (Suc i)) x - s (r i) x))* **if** $g' x \neq \infty$ **for** x **using** *that unfolding g'-def* **by** (*intro summable-suminf-not-top*) *fastforce* +

ultimately have *ae-summable: AE x in M. summable (λi. s (r (Suc i)) x - s (r i) x)* **using** *summable-norm-cancel* **unfolding** *dist-norm* **by** *force*

{
 fix x **assume** *summable (λi. s (r (Suc i)) x - s (r i) x)*
 hence $(\lambda n. \sum i < n. s (r (Suc i)) x - s (r i) x) \longrightarrow (\sum i. s (r (Suc i)) x - s (r i) x)$ **using** *summable-LIMSEQ* **by** *blast*
 moreover have $(\lambda n. (\sum i < n. s (r (Suc i)) x - s (r i) x)) = (\lambda n. s (r n) x - s (r 0) x)$ **using** *sum-lessThan-telescope* **by** *fastforce*
 ultimately have $(\lambda n. s (r n) x - s (r 0) x) \longrightarrow (\sum i. s (r (Suc i)) x - s (r i) x)$ **by** *argo*
 hence $(\lambda n. s (r n) x - s (r 0) x + s (r 0) x) \longrightarrow (\sum i. s (r (Suc i)) x - s (r i) x) + s (r 0) x$ **by** (*intro isCont-tendsto-compose[of - λz. z + s (r 0) x], auto*)
 hence *Cauchy (λn. s (r n) x)* **by** (*simp add: LIMSEQ-imp-Cauchy*)
}
hence *AE x in M. Cauchy (λi. s (r i) x)* **using** *ae-summable* **by** *fast*
thus *?thesis* **by** (*rule that[OF strict-mono(1)]*)
qed

3.2 Totally Ordered Banach Spaces

lemma *integrable-max[simp, intro]:*

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology}\}$

assumes *fg[measurable]: integrable M f integrable M g*

shows *integrable M (λx. max (f x) (g x))*

proof (*rule Bochner-Integration.integrable-bound*)

{
 fix $x y :: 'b$
 have $\text{norm } (\max x y) \leq \max (\text{norm } x) (\text{norm } y)$ **by** *linarith*
 also have $\dots \leq \text{norm } x + \text{norm } y$ **by** *simp*
 finally have $\text{norm } (\max x y) \leq \text{norm } (\text{norm } x + \text{norm } y)$ **by** *simp*
}
thus *AE x in M. norm (max (f x) (g x)) ≤ norm (norm (f x) + norm (g x))* **by** *simp*
qed (*auto intro: Bochner-Integration.integrable-add[OF integrable-norm[OF fg(1)]]*)

integrable-norm[*OF fg*(2)])])

lemma *integrable-min*[*simp*, *intro*]:

fixes *f* :: 'a \Rightarrow 'b :: {*second-countable-topology*, *banach*, *linorder-topology*}

assumes [*measurable*]: *integrable M f integrable M g*

shows *integrable M* ($\lambda x. \min (f x) (g x)$)

proof –

have *norm* ($\min (f x) (g x)$) \leq *norm* (*f x*) \vee *norm* ($\min (f x) (g x)$) \leq *norm* (*g x*) **for** *x* **by** *linarith*

thus ?thesis **by** (*intro integrable-bound*[*OF integrable-max*[*OF integrable-norm*(1,1), *OF assms*]], *auto*)

qed

lemma *integral-nonneg-AE-banach*:

fixes *f* :: 'a \Rightarrow 'b :: {*second-countable-topology*, *banach*, *linorder-topology*, *ordered-real-vector*}

assumes [*measurable*]: *f* \in *borel-measurable M* **and** *nonneg*: *AE x in M. 0 \leq f x*

shows $0 \leq \text{integral}^L M f$

proof *cases*

assume *integrable*: *integrable M f*

hence *max*: ($\lambda x. \max 0 (f x)$) \in *borel-measurable M* $\wedge x. 0 \leq \max 0 (f x)$ *integrable M* ($\lambda x. \max 0 (f x)$) **by** *auto*

hence $0 \leq \text{integral}^L M (\lambda x. \max 0 (f x))$

proof –

obtain *s* **where** *: $\wedge i. \text{simple-function } M (s i)$

$\wedge i. \text{emeasure } M \{y \in \text{space } M. s i y \neq 0\} \neq \infty$

$\wedge x. x \in \text{space } M \implies (\lambda i. s i x) \longrightarrow f x$

$\wedge x i. x \in \text{space } M \implies \text{norm } (s i x) \leq 2 * \text{norm } (f x)$ **using**

integrable-implies-simple-function-sequence[*OF integrable*] **by** *blast*

have *simple*: $\wedge i. \text{simple-function } M (\lambda x. \max 0 (s i x))$ **using** * **by** *fast*

have $\wedge i. \{y \in \text{space } M. \max 0 (s i y) \neq 0\} \subseteq \{y \in \text{space } M. s i y \neq 0\}$

unfolding *max-def* **by** *force*

moreover **have** $\wedge i. \{y \in \text{space } M. s i y \neq 0\} \in \text{sets } M$ **using** * **by** *measurable*

ultimately **have** $\wedge i. \text{emeasure } M \{y \in \text{space } M. \max 0 (s i y) \neq 0\} \leq \text{emeasure } M \{y \in \text{space } M. s i y \neq 0\}$ **by** (*rule emeasure-mono*)

hence **: $\wedge i. \text{emeasure } M \{y \in \text{space } M. \max 0 (s i y) \neq 0\} \neq \infty$ **using** *(2) **by** (*auto intro: order.strict-trans1 simp add: top.not-eq-extremum*)

have $\wedge x. x \in \text{space } M \implies (\lambda i. \max 0 (s i x)) \longrightarrow \max 0 (f x)$ **using** *(3) *tendsto-max* **by** *blast*

moreover **have** $\wedge x i. x \in \text{space } M \implies \text{norm } (\max 0 (s i x)) \leq \text{norm } (2 *_{\mathbb{R}} f x)$ **using** *(4) **unfolding** *max-def* **by** *auto*

ultimately **have** *tendsto*: $(\lambda i. \text{integral}^L M (\lambda x. \max 0 (s i x))) \longrightarrow \text{integral}^L M (\lambda x. \max 0 (f x))$

using *borel-measurable-simple-function simple integrable* **by** (*intro integral-dominated-convergence*[*OF max*(1) - *integrable-norm*[*OF integrable-scaleR-right*], *of - 2 f*], *blast+*)

{

fix *h* :: 'a \Rightarrow 'b :: {*second-countable-topology*, *banach*, *linorder-topology*, *ordered-real-vector*}


```

    assume simple-function M h emeasure M {y ∈ space M. h y ≠ 0} ≠ ∞ ∧ x.
x ∈ space M → h x ≥ 0
    hence *: integralL M h ≥ 0
    proof (induct rule: integrable-simple-function-induct-nn)
      case (cong f g)
      then show ?case using Bochner-Integration.integral-cong by force
    next
      case (indicator A y)
      hence A ≠ {} ⇒ y ≥ 0 using sets.sets-into-space by fastforce
      then show ?case using indicator by (cases A = {}, auto simp add:
scaleR-nonneg-nonneg)
    next
      case (add f g)
      then show ?case by (simp add: integrable-simple-function)
    qed
  }
  thus ?thesis using LIMSEQ-le-const[OF tendsto, of 0] ** simple by fastforce
qed
also have ... = integralL M f using nonneg by (auto intro: integral-cong-AE)
finally show ?thesis .
qed (simp add: not-integrable-integral-eq)

```

lemma *integral-mono-AE-banach*:

```

  fixes f g :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
dered-real-vector}
  assumes integrable M f integrable M g AE x in M. f x ≤ g x
  shows integralL M f ≤ integralL M g
  using integral-nonneg-AE-banach[of λx. g x - f x] assms Bochner-Integration.integral-diff[OF
assms(1,2)] by force

```

lemma *integral-mono-banach*:

```

  fixes f g :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
dered-real-vector}
  assumes integrable M f integrable M g ∧ x. x ∈ space M ⇒ f x ≤ g x
  shows integralL M f ≤ integralL M g
  using integral-mono-AE-banach assms by blast

```

3.3 Integrability and Measurability of the Diameter

context

```

  fixes s :: nat ⇒ 'a ⇒ 'b :: {second-countable-topology, banach} and M
  assumes bounded: ∧x. x ∈ space M ⇒ bounded (range (λi. s i x))
begin

```

lemma *borel-measurable-diameter*:

```

  assumes [measurable]: ∧i. (s i) ∈ borel-measurable M
  shows (λx. diameter {s i x | i. n ≤ i}) ∈ borel-measurable M
proof -
  have {s i x | i. N ≤ i} ≠ {} for x N by blast

```

hence *diameter-SUP*: $\text{diameter } \{s \ i \ x \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s \ i \ x) (s \ j \ x)) \text{ for } x \ N \text{ unfolding diameter-def by (auto intro!: arg-cong[of - Sup])}$

have *case-prod dist* ‘ $(\{s \ i \ x \mid i. n \leq i\} \times \{s \ i \ x \mid i. n \leq i\}) = ((\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) \text{ ‘ } (\{n..\} \times \{n..\}))$ ’ **for** x **by** *fast*

hence *: $(\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\}) = (\lambda x. \text{Sup } ((\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) \text{ ‘ } (\{n..\} \times \{n..\})))$ **using** *diameter-SUP* **by** (*simp add: case-prod-beta*)

have *bounded* $((\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) \text{ ‘ } (\{n..\} \times \{n..\}))$ **if** $x \in \text{space } M$ **for** x **by** (*rule bounded-imp-dist-bounded[OF bounded, OF that]*)

hence *bdd*: *bdd-above* $((\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) \text{ ‘ } (\{n..\} \times \{n..\}))$ **if** $x \in \text{space } M$ **for** x **using** *that bounded-imp-bdd-above* **by** *presburger*

have *fst* $p \in \text{borel-measurable } M$ *snd* $p \in \text{borel-measurable } M$ **if** $p \in s \text{ ‘ } \{n..\} \times s \text{ ‘ } \{n..\}$ **for** p **using** *that* **by** *fastforce+*

hence $(\lambda x. \text{fst } p \ x - \text{snd } p \ x) \in \text{borel-measurable } M$ **if** $p \in s \text{ ‘ } \{n..\} \times s \text{ ‘ } \{n..\}$ **for** p **using** *that borel-measurable-diff* **by** *simp*

hence $(\lambda x. \text{case } p \text{ of } (f, g) \Rightarrow \text{dist } (f \ x) (g \ x)) \in \text{borel-measurable } M$ **if** $p \in s \text{ ‘ } \{n..\} \times s \text{ ‘ } \{n..\}$ **for** p **unfolding** *dist-norm* **using** *that* **by** *measurable*

moreover **have** *countable* $(s \text{ ‘ } \{n..\} \times s \text{ ‘ } \{n..\})$ **by** (*intro countable-SIGMA countable-image, auto*)

ultimately show *?thesis* **unfolding** * **by** (*auto intro!: borel-measurable-cSUP bdd*)
qed

lemma *integrable-bound-diameter*:

fixes $f :: 'a \Rightarrow \text{real}$

assumes *integrable* $M \ f$

and [*measurable*]: $\bigwedge i. (s \ i) \in \text{borel-measurable } M$

and $\bigwedge x \ i. x \in \text{space } M \implies \text{norm } (s \ i \ x) \leq f \ x$

shows *integrable* $M \ (\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\})$

proof –

have $\{s \ i \ x \mid i. N \leq i\} \neq \{\}$ **for** $x \ N$ **by** *blast*

hence *diameter-SUP*: $\text{diameter } \{s \ i \ x \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s \ i \ x) (s \ j \ x)) \text{ for } x \ N \text{ unfolding diameter-def by (auto intro!: arg-cong[of - Sup])}$

{

fix x **assume** $x: x \in \text{space } M$

let $?S = (\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) \text{ ‘ } (\{n..\} \times \{n..\})$

have *case-prod dist* ‘ $(\{s \ i \ x \mid i. n \leq i\} \times \{s \ i \ x \mid i. n \leq i\}) = (\lambda(i, j). \text{dist } (s \ i \ x) (s \ j \ x)) \text{ ‘ } (\{n..\} \times \{n..\})$ ’ **by** *fast*

hence *: $\text{diameter } \{s \ i \ x \mid i. n \leq i\} = \text{Sup } ?S$ **using** *diameter-SUP* **by** (*simp add: case-prod-beta*)

have *bounded* $?S$ **by** (*rule bounded-imp-dist-bounded[OF bounded[OF x]]*)

hence *Sup-S-nonneg:0* $\leq \text{Sup } ?S$ **by** (*auto intro!: cSup-upper2 x bounded-imp-bdd-above*)

have $\text{dist } (s \ i \ x) (s \ j \ x) \leq 2 * f \ x$ **for** $i \ j$ **by** (*intro dist-triangle2[THEN order-trans, of - 0]*) (*metis norm-conv-dist assms(3) x add-mono mult-2*)

hence $\forall c \in ?S. c \leq 2 * f x$ **by** force
 hence $Sup ?S \leq 2 * f x$ **by** (intro cSup-least, auto)
 hence $norm (Sup ?S) \leq 2 * norm (f x)$ **using** Sup-S-nonneg **by** auto
 also have $\dots = norm (2 *_R f x)$ **by** simp
 finally have $norm (diameter \{s \ i \ x \mid i. n \leq i\}) \leq norm (2 *_R f x)$ **unfolding**
 * .
 }
 hence $\forall x \in M. norm (diameter \{s \ i \ x \mid i. n \leq i\}) \leq norm (2 *_R f x)$ **by** blast
 thus $integrable \ M (\lambda x. diameter \{s \ i \ x \mid i. n \leq i\})$ **using** borel-measurable-diameter
by (intro Bochner-Integration.integrable-bound[OF assms(1)[THEN integrable-scaleR-right[of
 2]]], measurable)
 qed
 end

3.4 Auxiliary Lemmas for Set Integrals

lemma set-integral-scaleR-left:

assumes $A \in sets \ M \ c \neq 0 \implies integrable \ M \ f$
 shows $LINT \ t:A \mid M. f \ t *_R c = (LINT \ t:A \mid M. f \ t) *_R c$
unfolding set-lebesgue-integral-def
using integrable-mult-indicator[OF assms]
by (subst integral-scaleR-left[symmetric], auto)

lemma nn-set-integral-eq-set-integral:

assumes [measurable]: $integrable \ M \ f$
 and $\forall x \in A \in M. 0 \leq f x \ A \in sets \ M$
 shows $(\int^+ x \in A. f x \ \partial M) = (\int x \in A. f x \ \partial M)$

proof –

have $(\int^+ x. indicator \ A \ x *_R f x \ \partial M) = (\int x \in A. f x \ \partial M)$
unfolding set-lebesgue-integral-def **using** assms(2) **by** (intro nn-integral-eq-integral[of
 - $\lambda x. indicat-real \ A \ x *_R f x$], blast intro: assms integrable-mult-indicator, fastforce)
 moreover have $(\int^+ x. indicator \ A \ x *_R f x \ \partial M) = (\int^+ x \in A. f x \ \partial M)$ **by** (metis
 ennreal-0 indicator-simps(1) indicator-simps(2) mult.commute mult-1 mult-zero-left
 real-scaleR-def)
 ultimately show ?thesis **by** argo
 qed

lemma set-integral-restrict-space:

fixes $f :: 'a \Rightarrow 'b :: \{banach, second-countable-topology\}$
 assumes $\Omega \cap space \ M \in sets \ M$
 shows $set-lebesgue-integral (restrict-space \ M \ \Omega) \ A \ f = set-lebesgue-integral \ M \ A$
 ($\lambda x. indicator \ \Omega \ x *_R f x$)
unfolding set-lebesgue-integral-def
by (subst integral-restrict-space, auto intro!: integrable-mult-indicator assms simp:
 mult.commute)

lemma set-integral-const:

fixes $c :: 'b :: \{banach, second-countable-topology\}$
 assumes $A \in sets \ M \ emeasure \ M \ A \neq \infty$

shows *set-lebesgue-integral* $M\ A\ (\lambda\cdot.\ c) = \text{measure } M\ A\ *_R\ c$
unfolding *set-lebesgue-integral-def*
using *assms* **by** (*metis has-bochner-integral-indicator has-bochner-integral-integral-eq infinity-enreal-def less-top*)

lemma *set-integral-mono-banach*:

fixes $f\ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes *set-integrable* $M\ A\ f$ *set-integrable* $M\ A\ g$
 $\bigwedge x. x \in A \implies f\ x \leq g\ x$
shows $(\text{LINT } x:A|M. f\ x) \leq (\text{LINT } x:A|M. g\ x)$
using *assms* **unfolding** *set-integrable-def set-lebesgue-integral-def*
by (*auto intro: integral-mono-banach split: split-indicator*)

lemma *set-integral-mono-AE-banach*:

fixes $f\ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes *set-integrable* $M\ A\ f$ *set-integrable* $M\ A\ g$ *AE* $x \in A$ *in* $M. f\ x \leq g\ x$
shows *set-lebesgue-integral* $M\ A\ f \leq \text{set-lebesgue-integral } M\ A\ g$ **using** *assms*
unfolding *set-lebesgue-integral-def* **by** (*auto simp add: set-integrable-def intro!: integral-mono-AE-banach[of M $\lambda x. \text{indicator } A\ x *_R f\ x$ $\lambda x. \text{indicator } A\ x *_R g\ x$], simp add: indicator-def*)

3.5 Averaging Theorem

lemma *balls-countable-basis*:

obtains $D :: 'a :: \{\text{metric-space, second-countable-topology}\}$ *set*
where *topological-basis* (*case-prod* *ball* ' ($D \times (\mathbb{Q} \cap \{0 < ..\})$)))
and *countable* D
and $D \neq \{\}$
proof –
obtain $D :: 'a$ *set* **where** *dense-subset: countable* $D\ D \neq \{\}$ $\llbracket \text{open } U; U \neq \{\} \rrbracket$
 $\implies \exists y \in D. y \in U$ **for** U **using** *countable-dense-exists* **by** *blast*
have *topological-basis* (*case-prod* *ball* ' ($D \times (\mathbb{Q} \cap \{0 < ..\})$)))
proof (*intro topological-basis-iff[THEN iffD2], fast, clarify*)
fix U **and** $x :: 'a$ **assume** *asm: open* $U\ x \in U$
obtain e **where** $e: e > 0$ *ball* $x\ e \subseteq U$ **using** *asm openE* **by** *blast*
obtain y **where** $y: y \in D\ y \in \text{ball } x\ (e / 3)$ **using** *dense-subset(3)[OF open-ball, of x e / 3] centre-in-ball[THEN iffD2, OF divide-pos-pos[OF e(1), of 3]]* **by** *force*
obtain r **where** $r: r \in \mathbb{Q} \cap \{e/3 < .. < e/2\}$ **unfolding** *Rats-def* **using** *of-rat-dense[OF divide-strict-left-mono[OF - e(1)], of 2 3]* **by** *auto*

have $*$: $x \in \text{ball } y\ r$ **using** $r\ y$ **by** (*simp add: dist-commute*)
hence $\text{ball } y\ r \subseteq U$ **using** r **by** (*intro order-trans[OF - e(2)], simp, metric*)
moreover **have** $\text{ball } y\ r \in (\text{case-prod } \text{ball} ' (D \times (\mathbb{Q} \cap \{0 < ..\})))$ **using** $y(1)$
by *force*
ultimately show $\exists B' \in (\text{case-prod } \text{ball} ' (D \times (\mathbb{Q} \cap \{0 < ..\}))). x \in B' \wedge B' \subseteq U$ **using** $*$ **by** *meson*
qed

thus ?thesis using that dense-subset by blast
qed

context sigma-finite-measure
begin

lemma sigma-finite-measure-induct[case-names finite-measure, consumes 0]:

assumes $\bigwedge(N :: 'a \text{ measure}) \Omega. \text{finite-measure } N$
 $\implies N = \text{restrict-space } M \ \Omega$
 $\implies \Omega \in \text{sets } M$
 $\implies \text{emeasure } N \ \Omega \neq \infty$
 $\implies \text{emeasure } N \ \Omega \neq 0$
 $\implies \text{almost-everywhere } N \ Q$

and [measurable]: Measurable.pred M Q

shows almost-everywhere M Q

proof -

have *: almost-everywhere N Q if finite-measure N $N = \text{restrict-space } M \ \Omega$ $\Omega \in \text{sets } M$ $\text{emeasure } N \ \Omega \neq \infty$ for N Ω using that by (cases emeasure N $\Omega = 0$, auto intro: emeasure-0-AE assms(1))

obtain A :: nat \Rightarrow 'a set where A: range A $\subseteq \text{sets } M$ $(\bigcup i. A \ i) = \text{space } M$ and emeasure-finite: emeasure M (A i) $\neq \infty$ for i using sigma-finite by metis

note A(1)[measurable]

have space-restr: space (restrict-space M (A i)) = A i for i unfolding space-restrict-space by simp

{

fix i

have *: $\{x \in A \ i \cap \text{space } M. \ Q \ x\} = \{x \in \text{space } M. \ Q \ x\} \cap (A \ i)$ by fast

have Measurable.pred (restrict-space M (A i)) Q using A by (intro measurableI, auto simp add: space-restr intro!: sets-restrict-space-iff[THEN iffD2], measurable, auto)

}

note this[measurable]

{

fix i

have finite-measure (restrict-space M (A i)) using emeasure-finite by (intro finite-measureI, subst space-restr, subst emeasure-restrict-space, auto)

hence emeasure (restrict-space M (A i)) $\{x \in A \ i. \neg Q \ x\} = 0$ using emeasure-finite by (intro AE-iff-measurable[THEN iffD1, OF - - *], measurable, subst space-restr[symmetric], intro sets.top, auto simp add: emeasure-restrict-space)

hence emeasure M $\{x \in A \ i. \neg Q \ x\} = 0$ by (subst emeasure-restrict-space[symmetric], auto)

}

hence emeasure M $(\bigcup i. \{x \in A \ i. \neg Q \ x\}) = 0$ by (intro emeasure-UN-eq-0, auto)

moreover have $(\bigcup i. \{x \in A \ i. \neg Q \ x\}) = \{x \in \text{space } M. \neg Q \ x\}$ using A by auto

ultimately show ?thesis by (intro AE-iff-measurable[THEN iffD2], auto)

qed

The following lemma allows us to make statements about the behaviour of a function almost everywhere, depending on the value it takes on average.

lemma *averaging-theorem*:

fixes $f :: \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes $[\text{measurable}]: \text{integrable } M f$
and $\text{closed: closed } S$
and $\bigwedge A. A \in \text{sets } M \implies \text{measure } M A > 0 \implies (1 / \text{measure } M A) *_R \text{set-lebesgue-integral } M A f \in S$
shows $\text{AE } x \text{ in } M. f x \in S$
proof (*induct rule: sigma-finite-measure-induct*)
case (*finite-measure* $N \Omega$)

interpret *finite-measure* N **by** (*rule finite-measure*)

have $\text{integrable}[\text{measurable}]: \text{integrable } N f$ **using** *assms finite-measure* **by** (*auto simp: integrable-restrict-space integrable-mult-indicator*)

have *average*: $(1 / \text{Sigma-Algebra.measure } N A) *_R \text{set-lebesgue-integral } N A f \in S$ **if** $A \in \text{sets } N \text{ measure } N A > 0$ **for** A

proof –

have $*$: $A \in \text{sets } M$ **using** *that* **by** (*simp add: sets-restrict-space-iff finite-measure*)

have $A = A \cap \Omega$ **by** (*metis finite-measure(2,3) inf.orderE sets.sets-into-space space-restrict-space that(1)*)

hence $\text{set-lebesgue-integral } N A f = \text{set-lebesgue-integral } M A f$ **unfolding** *finite-measure* **by** (*subst set-integral-restrict-space, auto simp add: finite-measure set-lebesgue-integral-def indicator-inter-arith[symmetric]*)

moreover **have** $\text{measure } N A = \text{measure } M A$ **using** *that* **by** (*auto intro!: measure-restrict-space simp add: finite-measure sets-restrict-space-iff*)

ultimately **show** *?thesis* **using** *that * assms(3)* **by** *presburger*

qed

obtain $D :: 'b \text{ set}$ **where** *balls-basis: topological-basis (case-prod ball ‘ (D × (Q ∩ {0<..})))* **and** *countable-D: countable D* **using** *balls-countable-basis* **by** *blast*

have *countable-balls: countable (case-prod ball ‘ (D × (Q ∩ {0<..})))* **using** *countable-rat countable-D* **by** *blast*

obtain B **where** *B-balls: B ⊆ case-prod ball ‘ (D × (Q ∩ {0<..})) ∪ B = −S* **using** *topological-basis[THEN iffD1, OF balls-basis] open-Compl[OF assms(2)]* **by** *meson*

hence *countable-B: countable B* **using** *countable-balls countable-subset* **by** *fast*

define b **where** $b = \text{from-nat-into } (B \cup \{\{\}\})$

have $B \cup \{\{\}\} \neq \{\}$ **by** *simp*

have *range-b: range b = B ∪ {\{\}}* **using** *countable-B* **by** (*auto simp add: b-def intro!: range-from-nat-into*)

have *open-b: open (b i)* **for** i **unfolding** *b-def* **using** *B-balls open-ball from-nat-into[of B ∪ {\{\}} i]* **by** *force*

have *Union-range-b: ∪(range b) = −S* **using** *B-balls range-b* **by** *simp*

{

```

fix  $v\ r$  assume ball-in-Compl:  $\text{ball } v\ r \subseteq -S$ 
define  $A$  where  $A = f - ' \text{ball } v\ r \cap \text{space } N$ 
have dist-less:  $\text{dist } (f\ x)\ v < r$  if  $x \in A$  for  $x$  using that unfolding  $A\text{-def}$ 
image-def by (simp add: dist-commute)
hence  $AE\text{-less}$ :  $AE\ x \in A$  in  $N$ .  $\text{norm } (f\ x - v) < r$  by (auto simp add:
dist-norm)
have  $*$ :  $A \in \text{sets } N$  unfolding  $A\text{-def}$  by simp
have emeasure  $N\ A = 0$ 
proof -
{
  assume asm: emeasure  $N\ A > 0$ 
  hence measure-pos: measure  $N\ A > 0$  unfolding emeasure-eq-measure by
simp
  hence  $(1 / \text{measure } N\ A) *_R \text{set-lebesgue-integral } N\ A\ f - v = (1 / \text{measure } N\ A) *_R \text{set-lebesgue-integral } N\ A\ (\lambda x. f\ x - v)$  using integrable integrable-const * by
(subst set-integral-diff(2), auto simp add: set-integrable-def set-integral-const[OF *]
algebra-simps intro!: integrable-mult-indicator)
  moreover have  $\text{norm } (\int_{x \in A}. (f\ x - v) \partial N) \leq (\int_{x \in A}. \text{norm } (f\ x - v) \partial N)$  using  $*$  by (auto intro!: integral-norm-bound[of N \lambda x. indicator A x *_R (f x - v), THEN order-trans] integrable-mult-indicator integrable simp add: set-lebesgue-integral-def)
  ultimately have  $\text{norm } ((1 / \text{measure } N\ A) *_R \text{set-lebesgue-integral } N\ A\ f - v) \leq \text{set-lebesgue-integral } N\ A\ (\lambda x. \text{norm } (f\ x - v)) / \text{measure } N\ A$  using asm
by (auto intro: divide-right-mono)
  also have  $\dots < \text{set-lebesgue-integral } N\ A\ (\lambda x. r) / \text{measure } N\ A$ 
  unfolding set-lebesgue-integral-def
  using asm * integrable integrable-const AE-less measure-pos
by (intro divide-strict-right-mono integral-less-AE[of - - A] integrable-mult-indicator
(fastforce simp add: dist-less dist-norm indicator-def)+)
  also have  $\dots = r$  using  $*$  measure-pos by (simp add: set-integral-const)
  finally have  $\text{dist } ((1 / \text{measure } N\ A) *_R \text{set-lebesgue-integral } N\ A\ f) v < r$ 
by (subst dist-norm)
  hence False using average[OF * measure-pos] by (metis ComplD dist-commute in-mono mem-ball ball-in-Compl)
}
thus ?thesis by fastforce
qed
}
note  $*$  = this
{
  fix  $b'$  assume  $b' \in B$ 
  hence ball-subset-Compl:  $b' \subseteq -S$  and ball-radius-pos:  $\exists v \in D. \exists r > 0. b' = \text{ball } v\ r$  using B-balls by (blast, fast)
}
note  $**$  = this
hence emeasure  $N\ (f - ' b\ i \cap \text{space } N) = 0$  for  $i$  by (cases b i = {}, simp)
(metis UnE singletonD * range-b[THEN eq-refl, THEN range-subsetD])
hence emeasure  $N\ (\bigcup i. f - ' b\ i \cap \text{space } N) = 0$  using open-b by (intro emeasure-UN-eq-0 fastforce+)

```

moreover have $(\bigcup i. f - ' b \ i \cap \text{space } N) = f - ' (\bigcup (\text{range } b)) \cap \text{space } N$ **by**
blast
ultimately have $\text{emeasure } N (f - ' (-S) \cap \text{space } N) = 0$ **using** *Union-range-b*
by *argo*
hence $AE\ x\ \text{in } N. f\ x \notin -S$ **using** *open-Compl[OF assms(2)]* **by** *(intro AE-iff-measurable[THEN iffD2], auto)*
thus *?case* **by** *force*
qed *(simp add: pred-sets2[OF borel-closed] assms(2))*

lemma *density-zero*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes *integrable M f*
and *density-0*: $\bigwedge A. A \in \text{sets } M \implies \text{set-lebesgue-integral } M\ A\ f = 0$
shows $AE\ x\ \text{in } M. f\ x = 0$
using *averaging-theorem[OF assms(1), of {0}] assms(2)*
by *(simp add: scaleR-nonneg-nonneg)*

This lemma shows that densities are unique in Banach spaces.

lemma *density-unique-banach*:
fixes $f\ f' :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes *integrable M f integrable M f'*
and *density-eq*: $\bigwedge A. A \in \text{sets } M \implies \text{set-lebesgue-integral } M\ A\ f = \text{set-lebesgue-integral } M\ A\ f'$
shows $AE\ x\ \text{in } M. f\ x = f'\ x$
proof –
{
fix A **assume** $asm: A \in \text{sets } M$
hence $LINT\ x|M. \text{indicat-real } A\ x\ *_R (f\ x - f'\ x) = 0$ **using** *density-eq*
assms(1,2) **by** *(simp add: set-lebesgue-integral-def algebra-simps Bochner-Integration.integral-diff[OF integrable-mult-indicator(1,1)])*
}
thus *?thesis* **using** *density-zero[OF Bochner-Integration.integrable-diff[OF assms(1,2)]]*
by *(simp add: set-lebesgue-integral-def)*
qed

lemma *density-nonneg*:
fixes $f :: - \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes *integrable M f*
and $\bigwedge A. A \in \text{sets } M \implies \text{set-lebesgue-integral } M\ A\ f \geq 0$
shows $AE\ x\ \text{in } M. f\ x \geq 0$
using *averaging-theorem[OF assms(1), of {0..}, OF closed-atLeast] assms(2)*
by *(simp add: scaleR-nonneg-nonneg)*

corollary *integral-nonneg-AE-eq-0-iff-AE*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, ordered-real-vector}\}$
assumes $f[\text{measurable}]: \text{integrable } M\ f$ **and** *nonneg*: $AE\ x\ \text{in } M. 0 \leq f\ x$
shows $\text{integral}^L\ M\ f = 0 \longleftrightarrow (AE\ x\ \text{in } M. f\ x = 0)$
proof


```

assume *:  $\text{integral}^L M f = 0$ 
{
  fix A assume asm:  $A \in \text{sets } M$ 
  have  $0 \leq \text{integral}^L M (\lambda x. \text{indicator } A x *_R f x)$  using nonneg by (subst inte-
    gral-zero[of M, symmetric], intro integral-mono-AE-banach integrable-mult-indicator
    asm f integrable-zero, auto simp add: indicator-def)
  moreover have  $\dots \leq \text{integral}^L M f$  using nonneg by (intro integral-mono-AE-banach
    integrable-mult-indicator asm f, auto simp add: indicator-def)
  ultimately have set-lebesgue-integral M A f = 0 unfolding set-lebesgue-integral-def
using * by force
}
thus AE x in M. f x = 0 by (intro density-zero f, blast)
qed (auto simp add: integral-eq-zero-AE)

```

corollary *integral-eq-mono-AE-eq-AE*:

```

fixes f g :: 'a  $\Rightarrow$  'b :: {second-countable-topology, banach, linorder-topology, or-
  dered-real-vector}
assumes integrable M f integrable M g  $\text{integral}^L M f = \text{integral}^L M g$  AE x in
  M. f x  $\leq$  g x
shows AE x in M. f x = g x
proof –
  define h where h =  $(\lambda x. g x - f x)$ 
  have AE x in M. h x = 0 unfolding h-def using assms by (subst inte-
    gral-nonneg-AE-eq-0-iff-AE[symmetric]) auto
  then show ?thesis unfolding h-def by auto
qed

```

end

end

theory *Conditional-Expectation-Banach*

```

imports HOL-Probability.Conditional-Expectation HOL-Probability.Independent-Family
  Bochner-Integration-Supplement
begin

```

4 Conditional Expectation in Banach Spaces

definition *has-cond-exp* :: 'a measure \Rightarrow 'a measure \Rightarrow ('a \Rightarrow 'b) \Rightarrow ('a \Rightarrow 'b :: {real-normed-vector, second-countable-topology}) \Rightarrow bool **where**

has-cond-exp M F f g = $((\forall A \in \text{sets } F. (\int x \in A. f x \, \partial M) = (\int x \in A. g x \, \partial M))$

$\wedge \text{integrable } M f$
 $\wedge \text{integrable } M g$
 $\wedge g \in \text{borel-measurable } F)$

lemma *has-cond-expI'*:

assumes $\bigwedge A. A \in \text{sets } F \implies (\int x \in A. f x \, \partial M) = (\int x \in A. g x \, \partial M)$

$integrable\ M\ f$
 $integrable\ M\ g$
 $g \in borel-measurable\ F$
shows $has-cond-exp\ M\ F\ f\ g$
using *assms* **unfolding** $has-cond-exp-def$ **by** *simp*

lemma $has-cond-expD$:
assumes $has-cond-exp\ M\ F\ f\ g$
shows $\bigwedge A. A \in sets\ F \implies (\int x \in A. f\ x\ \partial M) = (\int x \in A. g\ x\ \partial M)$
 $integrable\ M\ f$
 $integrable\ M\ g$
 $g \in borel-measurable\ F$
using *assms* **unfolding** $has-cond-exp-def$ **by** *simp+*

definition $cond-exp :: 'a\ measure \Rightarrow 'a\ measure \Rightarrow ('a \Rightarrow 'b) \Rightarrow ('a \Rightarrow 'b :: \{banach, second-countable-topology\})$ **where**
 $cond-exp\ M\ F\ f = (if\ \exists g. has-cond-exp\ M\ F\ f\ g\ then\ (SOME\ g. has-cond-exp\ M\ F\ f\ g)\ else\ (\lambda -. 0))$

lemma $borel-measurable-cond-exp[measurable]: cond-exp\ M\ F\ f \in borel-measurable\ F$
by (*metis* $cond-exp-def\ someI\ has-cond-exp-def\ borel-measurable-const$)

lemma $integrable-cond-exp[intro]: integrable\ M\ (cond-exp\ M\ F\ f)$
by (*metis* $cond-exp-def\ has-cond-expD(3)\ integrable-zero\ someI$)

lemma $set-integrable-cond-exp[intro]$:
assumes $A \in sets\ M$
shows $set-integrable\ M\ A\ (cond-exp\ M\ F\ f)$ **using** $integrable-mult-indicator[OF\ assms\ integrable-cond-exp,\ of\ F\ f]$ **by** (*auto* *simp* *add*: $set-integrable-def\ intro!$: $integrable-mult-indicator[OF\ assms\ integrable-cond-exp]$)

lemma $has-cond-exp-self$:
assumes $integrable\ M\ f$
shows $has-cond-exp\ M\ (vimage-algebra\ (space\ M)\ f\ borel)\ f\ f$
using *assms* **by** (*auto* *intro!*: $has-cond-expI'\ measurable-vimage-algebra1$)

lemma $has-cond-exp-sets-cong$:
assumes $sets\ F = sets\ G$
shows $has-cond-exp\ M\ F = has-cond-exp\ M\ G$
using *assms* **unfolding** $has-cond-exp-def$ **by** *force*

lemma $cond-exp-sets-cong$:
assumes $sets\ F = sets\ G$
shows $\forall x\ in\ M. cond-exp\ M\ F\ f\ x = cond-exp\ M\ G\ f\ x$
by (*intro* $AE-I2$, *simp* *add*: $cond-exp-def\ has-cond-exp-sets-cong[OF\ assms,\ of\ M]$)

context *sigma-finite-subalgebra*
begin

lemma *borel-measurable-cond-exp'[measurable]: cond-exp M F f ∈ borel-measurable M*
by (*metis cond-exp-def someI has-cond-exp-def borel-measurable-const subalg measurable-from-subalg*)

lemma *cond-exp-null:*
assumes $\nexists g. \text{has-cond-exp } M F f g$
shows $\text{cond-exp } M F f = (\lambda \cdot. 0)$
unfolding *cond-exp-def* **using** *assms* **by** *argo*

We state the tower property of the conditional expectation in terms of the predicate *has-cond-exp*.

lemma *has-cond-exp-nested-subalg:*
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes *subalgebra G F has-cond-exp M F f h has-cond-exp M G f h'*
shows $\text{has-cond-exp } M F h' h$
by (*intro has-cond-expI'*) (*metis assms has-cond-expD in-mono subalgebra-def*)**+**

The following lemma shows that the conditional expectation is unique as an element of L1, given that it exists.

lemma *has-cond-exp-charact:*
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$
assumes $\text{has-cond-exp } M F f g$
shows $\text{has-cond-exp } M F f (\text{cond-exp } M F f)$
 $AE\ x\ in\ M. \text{cond-exp } M F f\ x = g\ x$

proof –

show *cond-exp: has-cond-exp M F f (cond-exp M F f)* **using** *assms someI cond-exp-def* **by** *metis*

let $?MF = \text{restr-to-subalg } M F$

interpret *sigma-finite-measure ?MF* **by** (*rule sigma-fin-subalg*)

{

fix A **assume** $A \in \text{sets } ?MF$

then have $[measurable]: A \in \text{sets } F$ **using** *sets-restr-to-subalg[OF subalg]* **by** *simp*

have $(\int x \in A. g\ x\ \partial ?MF) = (\int x \in A. g\ x\ \partial M)$ **using** *assms subalg* **by** (*auto simp add: integral-subalgebra2 set-lebesgue-integral-def dest!: has-cond-expD*)

also have $\dots = (\int x \in A. \text{cond-exp } M F f\ x\ \partial M)$ **using** *assms cond-exp* **by** (*simp add: has-cond-exp-def*)

also have $\dots = (\int x \in A. \text{cond-exp } M F f\ x\ \partial ?MF)$ **using** *subalg* **by** (*auto simp add: integral-subalgebra2 set-lebesgue-integral-def*)

finally have $(\int x \in A. g\ x\ \partial ?MF) = (\int x \in A. \text{cond-exp } M F f\ x\ \partial ?MF)$ **by** *simp*

}

hence $AE\ x\ in\ ?MF. \text{cond-exp } M F f\ x = g\ x$ **using** *cond-exp assms subalg* **by** (*intro density-unique-banach, auto dest: has-cond-expD intro!: integrable-in-subalg*)

then show $AE\ x\ in\ M.\ cond_exp\ M\ F\ f\ x = g\ x$ **using** $AE_restr_to_subalg[OF\ subalg]$ **by** $simp$
qed

corollary $cond_exp_charact$:

fixes $f :: 'a \Rightarrow 'b :: \{second_countable_topology, banach\}$
assumes $\bigwedge A. A \in sets\ F \implies (\int x \in A. f\ x\ \partial M) = (\int x \in A. g\ x\ \partial M)$
 $integrable\ M\ f$
 $integrable\ M\ g$
 $g \in borel_measurable\ F$
shows $AE\ x\ in\ M.\ cond_exp\ M\ F\ f\ x = g\ x$
by $(intro\ has_cond_exp_charact\ has_cond_expI'\ assms)\ auto$

corollary $cond_exp_F_meas[intro, simp]$:

fixes $f :: 'a \Rightarrow 'b :: \{second_countable_topology, banach\}$
assumes $integrable\ M\ f$
 $f \in borel_measurable\ F$
shows $AE\ x\ in\ M.\ cond_exp\ M\ F\ f\ x = f\ x$
by $(rule\ cond_exp_charact, auto\ intro: assms)$

Congruence

lemma $has_cond_exp_cong$:

assumes $integrable\ M\ f\ \bigwedge x. x \in space\ M \implies f\ x = g\ x$ $has_cond_exp\ M\ F\ g\ h$
shows $has_cond_exp\ M\ F\ f\ h$
proof $(intro\ has_cond_expI'[OF\ -\ assms(1)], goal_cases)$
 $case\ (1\ A)$
hence $set_lebesgue_integral\ M\ A\ f = set_lebesgue_integral\ M\ A\ g$ **by** $(intro\ set_lebesgue_integral_cong)$
 $(meson\ assms(2)\ subalg\ in_mono\ subalgebra_def\ sets.sets_into_space\ subalgebra_def\ subsetD)+$
then show $?case$ **using** $1\ assms(3)$ **by** $(simp\ add: has_cond_exp_def)$
qed $(auto\ simp\ add: has_cond_expD[OF\ assms(3)])$

lemma $cond_exp_cong$:

fixes $f :: 'a \Rightarrow 'b :: \{second_countable_topology, banach\}$
assumes $integrable\ M\ f\ integrable\ M\ g\ \bigwedge x. x \in space\ M \implies f\ x = g\ x$
shows $AE\ x\ in\ M.\ cond_exp\ M\ F\ f\ x = cond_exp\ M\ F\ g\ x$
proof $(cases\ \exists h. has_cond_exp\ M\ F\ f\ h)$
 $case\ True$
then obtain h **where** $h: has_cond_exp\ M\ F\ f\ h\ has_cond_exp\ M\ F\ g\ h$ **using** $has_cond_exp_cong\ assms$ **by** $metis$
show $?thesis$ **using** $h[THEN\ has_cond_exp_charact(2)]$ **by** $fastforce$
next
 $case\ False$
moreover have $\nexists h. has_cond_exp\ M\ F\ g\ h$ **using** $False\ has_cond_exp_cong\ assms$
by $auto$
ultimately show $?thesis$ **unfolding** $cond_exp_def$ **by** $auto$
qed

lemma $has_cond_exp_cong_AE$:

assumes *integrable M f AE x in M. f x = g x has-cond-exp M F g h*
shows *has-cond-exp M F f h*
using *assms(1,2) subalg subalgebra-def subset-iff*
by (*intro has-cond-expI'*, *subst set-lebesgue-integral-cong-AE[OF - assms(1)[THEN borel-measurable-integrable] borel-measurable-integrable(1)[OF has-cond-expD(2)[OF assms(3)]]]*)
(fast intro: has-cond-expD[OF assms(3)] integrable-cong-AE-imp[OF - - AE-symmetric])+

lemma *has-cond-exp-cong-AE'*:
assumes *h ∈ borel-measurable F AE x in M. h x = h' x has-cond-exp M F f h'*
shows *has-cond-exp M F f h*
using *assms(1, 2) subalg subalgebra-def subset-iff*
using *AE-restr-to-subalg2[OF subalg assms(2)] measurable-from-subalg*
by (*intro has-cond-expI'*, *subst set-lebesgue-integral-cong-AE[OF - measurable-from-subalg(1,1)[OF subalg], OF - assms(1) has-cond-expD(4)[OF assms(3)]]]*)
(fast intro: has-cond-expD[OF assms(3)] integrable-cong-AE-imp[OF - - AE-symmetric])+

lemma *cond-exp-cong-AE*:
fixes *f :: 'a ⇒ 'b::{second-countable-topology,banach}*
assumes *integrable M f integrable M g AE x in M. f x = g x*
shows *AE x in M. cond-exp M F f x = cond-exp M F g x*
proof (*cases ∃ h. has-cond-exp M F f h*)
case *True*
then obtain h where *h: has-cond-exp M F f h has-cond-exp M F g h* **using** *has-cond-exp-cong-AE assms* **by** (*metis (mono-tags, lifting) eventually-mono*)
show *?thesis* **using** *h[THEN has-cond-exp-charact(2)]* **by** *fastforce*
next
case *False*
moreover have *¬ h. has-cond-exp M F g h* **using** *False has-cond-exp-cong-AE assms* **by** *auto*
ultimately show *?thesis* **unfolding** *cond-exp-def* **by** *auto*
qed

lemma *has-cond-exp-real*:
fixes *f :: 'a ⇒ real*
assumes *integrable M f*
shows *has-cond-exp M F f (real-cond-exp M F f)*
by (*intro has-cond-expI'*, *auto intro!: real-cond-exp-intA assms*)

lemma *cond-exp-real[intro]*:
fixes *f :: 'a ⇒ real*
assumes *integrable M f*
shows *AE x in M. cond-exp M F f x = real-cond-exp M F f x*
using *has-cond-exp-charact has-cond-exp-real assms* **by** *blast*

lemma *cond-exp-cmult*:
fixes *f :: 'a ⇒ real*
assumes *integrable M f*
shows *AE x in M. cond-exp M F (λx. c * f x) x = c * cond-exp M F f x*

using *real-cond-exp-cmult*[*OF* *assms*(1), *of* *c*] *assms*(1)[*THEN* *cond-exp-real*]
assms(1)[*THEN* *integrable-mult-right*, *THEN* *cond-exp-real*, *of* *c*] **by** *fastforce*

4.1 Existence

Indicator functions

lemma *has-cond-exp-indicator*:

assumes $A \in \text{sets } M \text{ emeasure } M \ A < \infty$

shows *has-cond-exp* $M \ F \ (\lambda x. \text{indicat-real } A \ x \ *_R \ y) \ (\lambda x. \text{real-cond-exp } M \ F \ (\text{indicator } A) \ x \ *_R \ y)$

proof (*intro* *has-cond-expI'*, *goal-cases*)

case (1 *B*)

have $\int x \in B. (\text{indicat-real } A \ x \ *_R \ y) \ \partial M = (\int x \in B. \text{indicat-real } A \ x \ \partial M) \ *_R \ y$
using *assms* **by** (*intro* *set-integral-scaleR-left*, *meson* 1 *in-mono subalg subalgebra-def*, *blast*)

also have $\dots = (\int x \in B. \text{real-cond-exp } M \ F \ (\text{indicator } A) \ x \ \partial M) \ *_R \ y$ **using** 1
assms **by** (*subst* *real-cond-exp-intA*, *auto*)

also have $\dots = \int x \in B. (\text{real-cond-exp } M \ F \ (\text{indicator } A) \ x \ *_R \ y) \ \partial M$ **using**
assms **by** (*intro* *set-integral-scaleR-left*[*symmetric*], *meson* 1 *in-mono subalg subalgebra-def*, *blast*)

finally show *?case* .

next

case 2

then show *?case* **using** *integrable-scaleR-left* *integrable-real-indicator* *assms* **by**
blast

next

case 3

show *?case* **using** *assms* **by** (*intro* *integrable-scaleR-left*, *intro* *real-cond-exp-int*,
blast+)

next

case 4

then show *?case* **by** (*intro* *borel-measurable-scaleR*, *intro* *Conditional-Expectation.borel-measurable-cond-exp*,
simp)

qed

lemma *cond-exp-indicator*[*intro*]:

fixes $y :: 'b :: \{\text{second-countable-topology}, \text{banach}\}$

assumes [*measurable*]: $A \in \text{sets } M \text{ emeasure } M \ A < \infty$

shows $\forall x \in M. \text{cond-exp } M \ F \ (\lambda x. \text{indicat-real } A \ x \ *_R \ y) \ x = \text{cond-exp } M \ F \ (\text{indicator } A) \ x \ *_R \ y$

proof –

have $\forall x \in M. \text{cond-exp } M \ F \ (\lambda x. \text{indicat-real } A \ x \ *_R \ y) \ x = \text{real-cond-exp } M \ F \ (\text{indicator } A) \ x \ *_R \ y$ **using** *has-cond-exp-indicator*[*OF* *assms*] *has-cond-exp-charact*
by *blast*

thus *?thesis* **using** *cond-exp-real*[*OF* *integrable-real-indicator*, *OF* *assms*] **by** *fastforce*

qed

Addition

```

lemma has-cond-exp-add:
  fixes  $f\ g :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
  assumes has-cond-exp  $M\ F\ f\ f'$  has-cond-exp  $M\ F\ g\ g'$ 
  shows has-cond-exp  $M\ F\ (\lambda x. f\ x + g\ x)\ (\lambda x. f'\ x + g'\ x)$ 
proof (intro has-cond-expI', goal-cases)
  case (1  $A$ )
    have  $\int_{x \in A}. (f\ x + g\ x) \partial M = (\int_{x \in A}. f\ x\ \partial M) + (\int_{x \in A}. g\ x\ \partial M)$  using
    assms[THEN has-cond-expD(2)] subalg 1 by (intro set-integral-add(2), auto simp
    add: subalgebra-def set-integrable-def intro: integrable-mult-indicator)
    also have  $\dots = (\int_{x \in A}. f'\ x\ \partial M) + (\int_{x \in A}. g'\ x\ \partial M)$  using assms[THEN
    has-cond-expD(1)[OF - 1]] by argo
    also have  $\dots = \int_{x \in A}. (f'\ x + g'\ x) \partial M$  using assms[THEN has-cond-expD(3)]
    subalg 1 by (intro set-integral-add(2)[symmetric], auto simp add: subalgebra-def
    set-integrable-def intro: integrable-mult-indicator)
    finally show ?case .
  next
    case 2
    then show ?case by (metis Bochner-Integration.integrable-add assms has-cond-expD(2))
  next
    case 3
    then show ?case by (metis Bochner-Integration.integrable-add assms has-cond-expD(3))
  next
    case 4
    then show ?case using assms borel-measurable-add has-cond-expD(4) by blast
qed

```

```

lemma has-cond-exp-scaleR-right:
  fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
  assumes has-cond-exp  $M\ F\ f\ f'$ 
  shows has-cond-exp  $M\ F\ (\lambda x. c *_{\mathbb{R}} f\ x)\ (\lambda x. c *_{\mathbb{R}} f'\ x)$ 
  using has-cond-expD[OF assms] by (intro has-cond-expI', auto)

```

```

lemma cond-exp-scaleR-right:
  fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
  assumes integrable  $M\ f$ 
  shows  $\text{AE } x \text{ in } M. \text{cond-exp } M\ F\ (\lambda x. c *_{\mathbb{R}} f\ x)\ x = c *_{\mathbb{R}} \text{cond-exp } M\ F\ f\ x$ 
proof (cases  $\exists f'. \text{has-cond-exp } M\ F\ f\ f'$ )
  case True
    then show ?thesis using assms has-cond-exp-charact has-cond-exp-scaleR-right
    by metis
  next
    case False
    show ?thesis
    proof (cases  $c = 0$ )
    case True
      then show ?thesis by simp
    next
      case c-nonzero: False
      have  $\nexists f'. \text{has-cond-exp } M\ F\ (\lambda x. c *_{\mathbb{R}} f\ x)\ f'$ 

```

```

proof (standard, goal-cases)
  case 1
  then obtain  $f'$  where  $f'$ : has-cond-exp  $M F (\lambda x. c *_R f x)$   $f'$  by blast
  have has-cond-exp  $M F f (\lambda x. \text{inverse } c *_R f' x)$  using has-cond-expD[OF
 $f'$ ] divideR-right[OF c-nonzero] assms by (intro has-cond-expI', auto)
  then show ?case using False by blast
qed
then show ?thesis using cond-exp-null[OF False] cond-exp-null by force
qed
qed

```

```

lemma cond-exp-uminus:
  fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
  assumes integrable  $M f$ 
  shows  $\text{AE } x \text{ in } M. \text{cond-exp } M F (\lambda x. - f x) x = - \text{cond-exp } M F f x$ 
  using cond-exp-scaleR-right[OF assms, of -1] by force

```

```

corollary has-cond-exp-simple:
  fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$ 
  assumes simple-function  $M f$  emeasure  $M \{y \in \text{space } M. f y \neq 0\} \neq \infty$ 
  shows has-cond-exp  $M F f (\text{cond-exp } M F f)$ 
  using assms
proof (induction rule: integrable-simple-function-induct)
  case (cong  $f g$ )
  then show ?case using has-cond-exp-cong by (metis (no-types, opaque-lifting)
Bochner-Integration.integrable-cong has-cond-expD(2) has-cond-exp-charact(1))
next
  case (indicator  $A y$ )
  then show ?case using has-cond-exp-charact[OF has-cond-exp-indicator] by fast
next
  case (add  $u v$ )
  then show ?case using has-cond-exp-add has-cond-exp-charact(1) by blast
qed

```

```

lemma cond-exp-contraction-real:
  fixes  $f :: 'a \Rightarrow \text{real}$ 
  assumes integrable[measurable]: integrable  $M f$ 
  shows  $\text{AE } x \text{ in } M. \text{norm} (\text{cond-exp } M F f x) \leq \text{cond-exp } M F (\lambda x. \text{norm} (f x)) x$ 
proof–
  have int: integrable  $M (\lambda x. \text{norm} (f x))$  using assms by blast
  have *:  $\text{AE } x \text{ in } M. 0 \leq \text{cond-exp } M F (\lambda x. \text{norm} (f x)) x$  using cond-exp-real[THEN
AE-symmetric, OF integrable-norm[OF integrable]] real-cond-exp-ge-c[OF integrable-norm[OF
integrable], of 0] norm-ge-zero by fastforce
  have **:  $A \in \text{sets } F \implies \int x \in A. |f x| \partial M = \int x \in A. \text{real-cond-exp } M F (\lambda x. \text{norm} (f x)) x \partial M$  for  $A$  unfolding real-norm-def using assms integrable-abs
real-cond-exp-intA by blast

```

```

  have norm-int:  $A \in \text{sets } F \implies (\int x \in A. |f x| \partial M) = (\int^+ x \in A. |f x| \partial M)$  for  $A$ 
using assms by (intro nn-set-integral-eq-set-integral[symmetric], blast, fastforce)

```


(meson subalg subalgebra-def subsetD)

have $AE\ x\ in\ M. real-cond-exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x \geq 0$ **using** *int real-cond-exp-ge-c*
by *force*

hence *cond-exp-norm-int*: $A \in sets\ F \implies (\int^{+} x \in A. real-cond-exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x\ \partial M) = (\int^{+} x \in A. real-cond-exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x\ \partial M)$ **for** A **using**
assms **by** (*intro nn-set-integral-eq-set-integral[symmetric]*, *blast*, *fastforce*) (*meson subalg subalgebra-def subsetD*)

have $A \in sets\ F \implies \int^{+} x \in A. |f\ x|\partial M = \int^{+} x \in A. real-cond-exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x\ \partial M$ **for** A **using** *** norm-int cond-exp-norm-int* **by** (*auto simp add: nn-integral-set-ennreal*)

moreover **have** $(\lambda x. ennreal\ |f\ x|) \in borel-measurable\ M$ **by** *measurable*

moreover **have** $(\lambda x. ennreal\ (real-cond-exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x)) \in borel-measurable\ F$ **by** *measurable*

ultimately **have** $AE\ x\ in\ M. nn-cond-exp\ M\ F\ (\lambda x. ennreal\ |f\ x|)\ x = real-cond-exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x$ **by** (*intro nn-cond-exp-charact[THEN AE-symmetric]*, *auto*)

hence $AE\ x\ in\ M. nn-cond-exp\ M\ F\ (\lambda x. ennreal\ |f\ x|)\ x \leq cond-exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x$ **using** *cond-exp-real[OF int]* **by** *force*

moreover **have** $AE\ x\ in\ M. |real-cond-exp\ M\ F\ f\ x| = norm\ (cond-exp\ M\ F\ f\ x)$

unfolding *real-norm-def* **using** *cond-exp-real[OF assms]* *** by** *force*

ultimately **have** $AE\ x\ in\ M. ennreal\ (norm\ (cond-exp\ M\ F\ f\ x)) \leq cond-exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x$ **using** *real-cond-exp-abs[OF assms[THEN borel-measurable-integrable]]*
by *fastforce*

hence $AE\ x\ in\ M. enn2real\ (ennreal\ (norm\ (cond-exp\ M\ F\ f\ x))) \leq enn2real\ (cond-exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x)$ **using** *ennreal-le-iff2* **by** *force*

thus *?thesis* **using** *** **by** *fastforce*

qed

lemma *cond-exp-contraction-simple*:

fixes $f :: 'a \Rightarrow 'b :: \{second-countable-topology, banach\}$

assumes *simple-function* $M\ f\ emeasure\ M\ \{y \in space\ M. f\ y \neq 0\} \neq \infty$

shows $AE\ x\ in\ M. norm\ (cond-exp\ M\ F\ f\ x) \leq cond-exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x$

using *assms*

proof (*induction rule: integrable-simple-function-induct*)

case (*cong f g*)

hence *ae*: $AE\ x\ in\ M. f\ x = g\ x$ **by** *blast*

hence $AE\ x\ in\ M. cond-exp\ M\ F\ f\ x = cond-exp\ M\ F\ g\ x$ **using** *cong has-cond-exp-simple*
by (*subst cond-exp-cong-AE*) (*auto intro!: has-cond-expD(2)*)

hence $AE\ x\ in\ M. norm\ (cond-exp\ M\ F\ f\ x) = norm\ (cond-exp\ M\ F\ g\ x)$ **by** *force*

moreover **have** $AE\ x\ in\ M. cond-exp\ M\ F\ (\lambda x. norm\ (f\ x))\ x = cond-exp\ M\ F\ (\lambda x. norm\ (g\ x))\ x$ **using** *ae cong has-cond-exp-simple* **by** (*subst cond-exp-cong-AE*)
(auto dest: has-cond-expD)

ultimately **show** *?case* **using** *cong(6)* **by** *fastforce*

next

case (*indicator A y*)

hence $AE\ x\ in\ M. cond-exp\ M\ F\ (\lambda a. indicator\ A\ a\ *_R\ y)\ x = cond-exp\ M\ F$

(indicator A) $x *_R y$ **by** blast

hence *: $AE\ x\ in\ M.\ norm\ (cond\text{-}exp\ M\ F\ (\lambda a.\ indicat\text{-}real\ A\ a *_R y)\ x) \leq norm\ y$
 $*\ cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (indicat\text{-}real\ A\ x))\ x$ **using** $cond\text{-}exp\text{-}contraction\text{-}real[OF\ integrable\text{-}real\text{-}indicator,\ OF\ indicator]$ **by** fastforce

have $AE\ x\ in\ M.\ norm\ y * cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (indicat\text{-}real\ A\ x))\ x = norm\ y$
 $*\ real\text{-}cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (indicat\text{-}real\ A\ x))\ x$ **using** $cond\text{-}exp\text{-}real[OF\ integrable\text{-}real\text{-}indicator,\ OF\ indicator]$ **by** fastforce

moreover **have** $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ y * norm\ (indicat\text{-}real\ A\ x))\ x =$
 $real\text{-}cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ y * norm\ (indicat\text{-}real\ A\ x))\ x$ **using**
 $indicator$ **by** (intro $cond\text{-}exp\text{-}real$, auto)

ultimately **have** $AE\ x\ in\ M.\ norm\ y * cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (indicat\text{-}real\ A\ x))\ x =$
 $cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ y * norm\ (indicat\text{-}real\ A\ x))\ x$ **using** $real\text{-}cond\text{-}exp\text{-}cmult[of\ \lambda x.\ norm\ (indicat\text{-}real\ A\ x)\ norm\ y]$ $indicator$ **by** fastforce

moreover **have** $(\lambda x.\ norm\ y * norm\ (indicat\text{-}real\ A\ x)) = (\lambda x.\ norm\ (indicat\text{-}real\ A\ x *_R y))$
by force

ultimately **show** ?case **using** * **by** force

next

case (add $u\ v$)

have $AE\ x\ in\ M.\ norm\ (cond\text{-}exp\ M\ F\ (\lambda a.\ u\ a + v\ a)\ x) = norm\ (cond\text{-}exp\ M\ F\ u\ x +$
 $cond\text{-}exp\ M\ F\ v\ x)$ **using** $has\text{-}cond\text{-}exp\text{-}character(2)[OF\ has\text{-}cond\text{-}exp\text{-}add,\ OF\ has\text{-}cond\text{-}exp\text{-}simple(1,1),\ OF\ add(1,2,3,4)]$ **by** fastforce

moreover **have** $AE\ x\ in\ M.\ norm\ (cond\text{-}exp\ M\ F\ u\ x + cond\text{-}exp\ M\ F\ v\ x) \leq$
 $norm\ (cond\text{-}exp\ M\ F\ u\ x) + norm\ (cond\text{-}exp\ M\ F\ v\ x)$ **using** $norm\text{-}triangle\text{-}ineq$
by blast

moreover **have** $AE\ x\ in\ M.\ norm\ (cond\text{-}exp\ M\ F\ u\ x) + norm\ (cond\text{-}exp\ M\ F\ v\ x) \leq$
 $cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (u\ x))\ x + cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (v\ x))\ x$ **using**
 $add(6,7)$ **by** fastforce

moreover **have** $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (u\ x))\ x + cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (v\ x))\ x =$
 $cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (u\ x) + norm\ (v\ x))\ x$ **using** $integrable\text{-}simple\text{-}function[OF\ add(1,2)]$
 $integrable\text{-}simple\text{-}function[OF\ add(3,4)]$ **by** (intro $has\text{-}cond\text{-}exp\text{-}character(2)[OF\ has\text{-}cond\text{-}exp\text{-}add[OF\ has\text{-}cond\text{-}exp\text{-}character(1,1)],$
 $THEN\ AE\text{-}symmetric]$, auto intro: $has\text{-}cond\text{-}exp\text{-}real$)

moreover **have** $AE\ x\ in\ M.\ cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (u\ x) + norm\ (v\ x))\ x =$
 $cond\text{-}exp\ M\ F\ (\lambda x.\ norm\ (u\ x + v\ x))\ x$ **using** $add(5)\ integrable\text{-}simple\text{-}function[OF\ add(1,2)]$
 $integrable\text{-}simple\text{-}function[OF\ add(3,4)]$ **by** (intro $cond\text{-}exp\text{-}cong$, auto)

ultimately **show** ?case **by** force

qed

lemma $has\text{-}cond\text{-}exp\text{-}simple\text{-}lim$:

fixes $f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology,\ banach\}$

assumes $integrable[measurable]: integrable\ M\ f$

and $\bigwedge i.\ simple\text{-}function\ M\ (s\ i)$

and $\bigwedge i.\ emeasure\ M\ \{y \in space\ M.\ s\ i\ y \neq 0\} \neq \infty$

and $\bigwedge x.\ x \in space\ M \implies (\lambda i.\ s\ i\ x) \longrightarrow f\ x$

and $\bigwedge x\ i.\ x \in space\ M \implies norm\ (s\ i\ x) \leq 2 * norm\ (f\ x)$

obtains r

where $strict\text{-}mono\ r\ has\text{-}cond\text{-}exp\ M\ F\ f\ (\lambda x.\ lim\ (\lambda i.\ cond\text{-}exp\ M\ F\ (s\ (r\ i))\ x))$

$AE\ x\ in\ M.\ convergent\ (\lambda i.\ cond_exp\ M\ F\ (s\ (r\ i))\ x)$
proof –
 have $[measurable]: (s\ i) \in borel_measurable\ M$ **for** i **using** $assms(2)$ **by** $(simp\ add:\ borel_measurable_simple_function)$
 have $integrable_s: integrable\ M\ (\lambda x.\ s\ i\ x)$ **for** i **using** $assms\ integrable_simple_function$
by $blast$
 have $integrable_4f: integrable\ M\ (\lambda x.\ 4 * norm\ (f\ x))$ **using** $assms(1)$ **by** $simp$
 have $integrable_2f: integrable\ M\ (\lambda x.\ 2 * norm\ (f\ x))$ **using** $assms(1)$ **by** $simp$
 have $integrable_2_cond_exp_norm_f: integrable\ M\ (\lambda x.\ 2 * cond_exp\ M\ F\ (\lambda x.\ norm\ (f\ x))\ x)$ **by** $fast$

 have $emeasure\ M\ \{y \in space\ M.\ s\ i\ y - s\ j\ y \neq 0\} \leq emeasure\ M\ \{y \in space\ M.\ s\ i\ y \neq 0\} + emeasure\ M\ \{y \in space\ M.\ s\ j\ y \neq 0\}$ **for** $i\ j$ **using** $simple_functionD(2)[OF\ assms(2)]$ **by** $(intro\ order_trans[OF\ emeasure_mono\ emeasure_subadditive],\ auto)$
 hence $fin_sup: emeasure\ M\ \{y \in space\ M.\ s\ i\ y - s\ j\ y \neq 0\} \neq \infty$ **for** $i\ j$ **using** $assms(3)$ **by** $(metis\ (mono_tags)\ ennreal_add_eq_top\ linorder_not_less\ top.not_eq_extremum\ infinity_ennreal_def)$

 have $emeasure\ M\ \{y \in space\ M.\ norm\ (s\ i\ y - s\ j\ y) \neq 0\} \leq emeasure\ M\ \{y \in space\ M.\ s\ i\ y \neq 0\} + emeasure\ M\ \{y \in space\ M.\ s\ j\ y \neq 0\}$ **for** $i\ j$ **using** $simple_functionD(2)[OF\ assms(2)]$ **by** $(intro\ order_trans[OF\ emeasure_mono\ emeasure_subadditive],\ auto)$
 hence $fin_sup_norm: emeasure\ M\ \{y \in space\ M.\ norm\ (s\ i\ y - s\ j\ y) \neq 0\} \neq \infty$ **for** $i\ j$ **using** $assms(3)$ **by** $(metis\ (mono_tags)\ ennreal_add_eq_top\ linorder_not_less\ top.not_eq_extremum\ infinity_ennreal_def)$

 have $Cauchy: Cauchy\ (\lambda n.\ s\ n\ x)$ **if** $x \in space\ M$ **for** x **using** $assms(4)$ $LIM_SEQ_imp_Cauchy$ **that** **by** $blast$
 hence $bounded_range_s: bounded\ (range\ (\lambda n.\ s\ n\ x))$ **if** $x \in space\ M$ **for** x **using** $that\ cauchy_imp_bounded$ **by** $fast$

 have $AE\ x\ in\ M.\ (\lambda n.\ diameter\ \{s\ i\ x \mid i.\ n \leq i\}) \longrightarrow 0$ **using** $Cauchy\ cauchy_iff_diameter_tends_to_zero_and_bounded$ **by** $fast$
 moreover have $(\lambda x.\ diameter\ \{s\ i\ x \mid i.\ n \leq i\}) \in borel_measurable\ M$ **for** n **using** $bounded_range_s\ borel_measurable_diameter$ **by** $measurable$
 moreover have $AE\ x\ in\ M.\ norm\ (diameter\ \{s\ i\ x \mid i.\ n \leq i\}) \leq 4 * norm\ (f\ x)$ **for** n
proof –
 {
 fix x **assume** $x: x \in space\ M$
 have $diameter\ \{s\ i\ x \mid i.\ n \leq i\} \leq 2 * norm\ (f\ x) + 2 * norm\ (f\ x)$
by $(intro\ diameter_le,\ blast,\ subst\ dist_norm[symmetric],\ intro\ dist_triangle3[THEN\ order_trans,\ of\ 0],\ intro\ add_mono)\ (auto\ intro:\ assms(5)[OF\ x])$
 hence $norm\ (diameter\ \{s\ i\ x \mid i.\ n \leq i\}) \leq 4 * norm\ (f\ x)$ **using** $diameter_ge_0[OF\ bounded_subset[OF\ bounded_range_s],\ OF\ x,\ of\ \{s\ i\ x \mid i.\ n \leq i\}]$ **by** $force$
 }
thus $?thesis$ **by** $fast$

qed

ultimately have *diameter-tendsto-zero*: $(\lambda n. \text{LINT } x | M. \text{diameter } \{s \ i \ x \mid i. n \leq i\}) \longrightarrow 0$ **by** (intro *integral-dominated-convergence*[OF *borel-measurable-const*[of 0] - *integrable-4f*, *simplified*]) (fast+)

have *diameter-integrable*: *integrable* M $(\lambda x. \text{diameter } \{s \ i \ x \mid i. n \leq i\})$ **for** n **using** *assms*(1,5)

by (intro *integrable-bound-diameter*[OF *bounded-range-s integrable-2f*], *auto*)

have *dist-integrable*: *integrable* M $(\lambda x. \text{dist } (s \ i \ x) (s \ j \ x))$ **for** $i \ j$ **using** *assms*(5)

dist-triangle3[of $s \ i \ - \ 0$, THEN *order-trans*, OF *add-mono*, of $- \ 2 * \text{norm } (f \ -)$]

by (intro *Bochner-Integration.integrable-bound*[OF *integrable-4f*]) *fastforce*+

have $\exists N. \forall i \geq N. \forall j \geq N. \text{LINT } x | M. \text{norm } (\text{cond-exp } M \ F \ (s \ i) \ x - \text{cond-exp } M \ F \ (s \ j) \ x) < e$ **if** $e\text{-pos}$: $e > 0$ **for** e

proof -

obtain N **where** *: *LINT* $x | M. \text{diameter } \{s \ i \ x \mid i. n \leq i\} < e$ **if** $n \geq N$ **for** n **using** *that order-tendsto-iff*[THEN *iffD1*, OF *diameter-tendsto-zero*, *unfolded eventually-sequentially*] $e\text{-pos}$ **by** *presburger*

{

fix $i \ j \ x$ **assume** *asm*: $i \geq N \ j \geq N \ x \in \text{space } M$

have *case-prod dist* ' $(\{s \ i \ x \mid i. N \leq i\} \times \{s \ i \ x \mid i. N \leq i\}) = \text{case-prod } (\lambda i \ j. \text{dist } (s \ i \ x) (s \ j \ x)) \text{ ' } (\{N..\} \times \{N..\})$ **by** *fast*

hence *diameter* $\{s \ i \ x \mid i. N \leq i\} = (\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s \ i \ x) (s \ j \ x))$ **unfolding** *diameter-def* **by** *auto*

moreover have $(\text{SUP } (i, j) \in \{N..\} \times \{N..\}. \text{dist } (s \ i \ x) (s \ j \ x)) \geq \text{dist } (s \ i \ x) (s \ j \ x)$ **using** *asm bounded-imp-bdd-above*[OF *bounded-imp-dist-bounded*, OF *bounded-range-s*] **by** (intro *cSup-upper*, *auto*)

ultimately have *diameter* $\{s \ i \ x \mid i. N \leq i\} \geq \text{dist } (s \ i \ x) (s \ j \ x)$ **by** *presburger*

}

hence *LINT* $x | M. \text{dist } (s \ i \ x) (s \ j \ x) < e$ **if** $i \geq N \ j \geq N$ **for** $i \ j$ **using** *that ** **by** (intro *integral-mono*[OF *dist-integrable diameter-integrable*, THEN *order.strict-trans1*], *blast*+))

moreover have *LINT* $x | M. \text{norm } (\text{cond-exp } M \ F \ (s \ i) \ x - \text{cond-exp } M \ F \ (s \ j) \ x) \leq \text{LINT } x | M. \text{dist } (s \ i \ x) (s \ j \ x)$ **for** $i \ j$

proof -

have *LINT* $x | M. \text{norm } (\text{cond-exp } M \ F \ (s \ i) \ x - \text{cond-exp } M \ F \ (s \ j) \ x) = \text{LINT } x | M. \text{norm } (\text{cond-exp } M \ F \ (s \ i) \ x + - \ 1 * \text{cond-exp } M \ F \ (s \ j) \ x)$ **unfolding** *dist-norm* **by** *simp*

also have ... = *LINT* $x | M. \text{norm } (\text{cond-exp } M \ F \ (\lambda x. s \ i \ x - s \ j \ x) \ x)$ **using** *has-cond-exp-charact*(2)[OF *has-cond-exp-add*[OF *has-cond-exp-scaleR-right*, OF *has-cond-exp-charact*(1,1), OF *has-cond-exp-simple*(1,1)[OF *assms*(2,3)]]], THEN *AE-symmetric*, of $i - 1 \ j$] **by** (intro *integral-cong-AE*) *force*+

also have ... $\leq \text{LINT } x | M. \text{cond-exp } M \ F \ (\lambda x. \text{norm } (s \ i \ x - s \ j \ x)) \ x$ **using** *cond-exp-contraction-simple*[OF *fin-sup*, of $i \ j$] *integrable-cond-exp assms*(2) **by** (intro *integral-mono-AE*, *fast*+))

also have ... = *LINT* $x | M. \text{norm } (s \ i \ x - s \ j \ x)$ **unfolding** *set-integral-space*(1)[OF *integrable-cond-exp, symmetric*] *set-integral-space*[OF *dist-integrable*[*unfolded dist-norm*],

symmetric] **by** (*intro has-cond-expD*(1)[*OF has-cond-exp-simple*[*OF - fin-sup-norm*], *symmetric*]) (*metis assms*(2) *simple-function-compose1 simple-function-diff*, *metis sets.top subalg subalgebra-def*)

finally show *?thesis unfolding dist-norm* .

qed

ultimately show *?thesis using order.strict-trans1* **by** *meson*

qed

then obtain *r* **where** *strict-mono-r: strict-mono r* **and** *AE-Cauchy: AE x in M. Cauchy (λi. cond-exp M F (s (r i)) x)* **by** (*rule cauchy-L1-AE-cauchy-subseq*[*OF integrable-cond-exp*], *auto*)

hence *ae-lim-cond-exp: AE x in M. (λn. cond-exp M F (s (r n)) x) ⟶ lim (λn. cond-exp M F (s (r n)) x)* **using** *Cauchy-convergent-iff convergent-LIMSEQ-iff* **by** *fastforce*

have *cond-exp-bounded: AE x in M. norm (cond-exp M F (s (r n)) x) ≤ cond-exp M F (λx. 2 * norm (f x)) x* **for** *n*

proof –

have *AE x in M. norm (cond-exp M F (s (r n)) x) ≤ cond-exp M F (λx. norm (s (r n) x)) x* **by** (*rule cond-exp-contraction-simple*[*OF assms*(2,3)])

moreover have *AE x in M. real-cond-exp M F (λx. norm (s (r n) x)) x ≤ real-cond-exp M F (λx. 2 * norm (f x)) x* **using** *integrable-s integrable-2f assms*(5) **by** (*intro real-cond-exp-mono*, *auto*)

ultimately show *?thesis using cond-exp-real*[*OF integrable-norm*, *OF integrable-s*, *of r n*] *cond-exp-real*[*OF integrable-2f*] **by** *force*

qed

have *lim-integrable: integrable M (λx. lim (λi. cond-exp M F (s (r i)) x))* **by** (*intro integrable-dominated-convergence*[*OF - borel-measurable-cond-exp' integrable-cond-exp ae-lim-cond-exp cond-exp-bounded*], *simp*)

{

fix *A* **assume** *A-in-sets-F: A ∈ sets F*

have *AE x in M. norm (indicator A x *_R cond-exp M F (s (r n)) x) ≤ cond-exp M F (λx. 2 * norm (f x)) x* **for** *n*

proof –

have *AE x in M. norm (indicator A x *_R cond-exp M F (s (r n)) x) ≤ norm (cond-exp M F (s (r n)) x)* **unfolding** *indicator-def* **by** *simp*

thus *?thesis using cond-exp-bounded*[*of n*] **by** *force*

qed

hence *lim-cond-exp-int: (λn. LINT x:A|M. cond-exp M F (s (r n)) x) ⟶ LINT x:A|M. lim (λn. cond-exp M F (s (r n)) x)*

using *ae-lim-cond-exp measurable-from-subalg*[*OF subalg borel-measurable-indicator, OF A-in-sets-F*] *cond-exp-bounded*

unfolding *set-lebesgue-integral-def*

by (*intro integral-dominated-convergence*[*OF borel-measurable-scaleR borel-measurable-scaleR integrable-cond-exp*]) (*fastforce simp add: tendsto-scaleR*) +

have *AE x in M. norm (indicator A x *_R s (r n) x) ≤ 2 * norm (f x)* **for** *n*

proof –

have *AE x in M. norm (indicator A x *_R s (r n) x) ≤ norm (s (r n) x)*

unfolding *indicator-def* **by** *simp*
thus *?thesis* **using** *assms(5)[of - r n]* **by** *fastforce*
qed
hence *lim-s-int*: $(\lambda n. \text{LINT } x:A|M. s (r n) x) \longrightarrow \text{LINT } x:A|M. f x$
using *measurable-from-subalg[OF subalg borel-measurable-indicator, OF A-in-sets-F]*
LIMSEQ-subseq-LIMSEQ[OF assms(4) strict-mono-r] *assms(5)*
unfolding *set-lebesgue-integral-def comp-def*
by (*intro integral-dominated-convergence[OF borel-measurable-scaleR borel-measurable-scaleR integrable-2f]*) (*fastforce simp add: tendsto-scaleR*)

have *LINT x:A|M. lim* ($\lambda n. \text{cond-exp } M F (s (r n)) x$) = *lim* ($\lambda n. \text{LINT } x:A|M. \text{cond-exp } M F (s (r n)) x$) **using** *limI[OF lim-cond-exp-int]* **by** *argo*
also have $\dots = \text{lim} (\lambda n. \text{LINT } x:A|M. s (r n) x)$ **using** *has-cond-expD(1)[OF has-cond-exp-simple[OF assms(2,3)] A-in-sets-F, symmetric]* **by** *presburger*
also have $\dots = \text{LINT } x:A|M. f x$ **using** *limI[OF lim-s-int]* **by** *argo*
finally have *LINT x:A|M. lim* ($\lambda n. \text{cond-exp } M F (s (r n)) x$) = *LINT x:A|M. f x* .
}
hence *has-cond-exp M F f* ($\lambda x. \text{lim} (\lambda i. \text{cond-exp } M F (s (r i)) x)$) **using** *assms(1) lim-integrable* **by** (*intro has-cond-expI', auto*)
thus *thesis* **using** *AE-Cauchy Cauchy-convergent strict-mono-r* **by** (*auto intro!: that*)
qed

corollary *has-cond-expI*:
fixes *f :: 'a \Rightarrow 'b::{second-countable-topology,banach}*
assumes *integrable M f*
shows *has-cond-exp M F f* (*cond-exp M F f*)
proof –
obtain *s* **where** *s-is*: $\bigwedge i. \text{simple-function } M (s i) \bigwedge i. \text{emeasure } M \{y \in \text{space } M. s i y \neq 0\} \neq \infty \bigwedge x. x \in \text{space } M \implies (\lambda i. s i x) \longrightarrow f x \bigwedge x i. x \in \text{space } M \implies \text{norm } (s i x) \leq 2 * \text{norm } (f x)$ **using** *integrable-implies-simple-function-sequence[OF assms]* **by** *blast*
show *?thesis* **using** *has-cond-exp-simple-lim[OF assms s-is]* *has-cond-exp-charact(1)*
by *metis*
qed

4.2 Properties

lemma *cond-exp-nested-subalg*:

fixes *f :: 'a \Rightarrow 'b::{second-countable-topology,banach}*
assumes *integrable M f subalgebra M G subalgebra G F*
shows *AE ξ in M. cond-exp M F f ξ = cond-exp M F (cond-exp M G f) ξ*
using *has-cond-expI assms sigma-finite-subalgebra-def* **by** (*auto intro!: has-cond-exp-nested-subalg[THEN has-cond-exp-charact(2), THEN AE-symmetric] sigma-finite-subalgebra.has-cond-expI[OF sigma-finite-subalgebra.intro[OF assms(2)]] nested-subalg-is-sigma-finite*)

lemma *cond-exp-set-integral*:

fixes *f :: 'a \Rightarrow 'b::{second-countable-topology,banach}*

assumes *integrable* $M f A \in \text{sets } F$
shows $(\int x \in A. f x \partial M) = (\int x \in A. \text{cond-exp } M F f x \partial M)$
using *has-cond-expD(1)*[*OF has-cond-expI*, *OF assms*] **by** *argo*

lemma *cond-exp-add*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f$ *integrable* $M g$
shows $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x + g x) x = \text{cond-exp } M F f x + \text{cond-exp } M F g x$
using *has-cond-exp-add*[*OF has-cond-expI(1,1)*, *OF assms*, *THEN has-cond-exp-charact(2)*]
.

lemma *cond-exp-diff*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f$ *integrable* $M g$
shows $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x - g x) x = \text{cond-exp } M F f x - \text{cond-exp } M F g x$
using *has-cond-exp-add*[*OF - has-cond-exp-scaleR-right*, *OF has-cond-expI(1,1)*, *OF assms*, *THEN has-cond-exp-charact(2)*, *of -1*] **by** *simp*

lemma *cond-exp-diff'*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f$ *integrable* $M g$
shows $AE x \text{ in } M. \text{cond-exp } M F (f - g) x = \text{cond-exp } M F f x - \text{cond-exp } M F g x$
unfolding *fun-diff-def* **using** *assms* **by** (*rule cond-exp-diff*)

lemma *cond-exp-scaleR-left*:
fixes $f :: 'a \Rightarrow \text{real}$
assumes *integrable* $M f$
shows $AE x \text{ in } M. \text{cond-exp } M F (\lambda x. f x *_R c) x = \text{cond-exp } M F f x *_R c$
using *cond-exp-set-integral*[*OF assms*] *subalg assms* **unfolding** *subalgebra-def*
by (*intro cond-exp-charact*,
subst set-integral-scaleR-left, *blast*, *intro assms*,
subst set-integral-scaleR-left, *blast*, *intro integrable-cond-exp*)
auto

lemma *cond-exp-contraction*:
fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}\}$
assumes *integrable* $M f$
shows $AE x \text{ in } M. \text{norm } (\text{cond-exp } M F f x) \leq \text{cond-exp } M F (\lambda x. \text{norm } (f x)) x$
proof –
obtain s **where** $s: \bigwedge i. \text{simple-function } M (s i) \bigwedge i. \text{emeasure } M \{y \in \text{space } M. s i y \neq 0\} \neq \infty \bigwedge x. x \in \text{space } M \implies (\lambda i. s i x) \longrightarrow f x \bigwedge i x. x \in \text{space } M$

$\implies \text{norm } (s \ i \ x) \leq 2 * \text{norm } (f \ x)$

by (*blast intro: integrable-implies-simple-function-sequence*[*OF assms*])

obtain r **where** r : *strict-mono* r **and** *has-cond-exp* $M \ F \ f \ (\lambda x. \lim (\lambda i. \text{cond-exp } M \ F \ (s \ (r \ i)) \ x)) \ AE \ x \text{ in } M. (\lambda i. \text{cond-exp } M \ F \ (s \ (r \ i)) \ x) \longrightarrow \lim (\lambda i. \text{cond-exp } M \ F \ (s \ (r \ i)) \ x)$

using *has-cond-exp-simple-lim*[*OF assms s*] **unfolding** *convergent-LIMSEQ-iff* **by** *blast*

hence $r\text{-tendsto}$: $AE \ x \text{ in } M. (\lambda i. \text{cond-exp } M \ F \ (s \ (r \ i)) \ x) \longrightarrow \text{cond-exp } M \ F \ f \ x$ **using** *has-cond-exp-charact*(2) **by** *force*

have norm-s-r : $\bigwedge i. \text{simple-function } M \ (\lambda x. \text{norm } (s \ (r \ i) \ x)) \bigwedge i. \text{emeasure } M \ \{y \in \text{space } M. \text{norm } (s \ (r \ i) \ y) \neq 0\} \neq \infty \bigwedge x. x \in \text{space } M \implies (\lambda i. \text{norm } (s \ (r \ i) \ x)) \longrightarrow \text{norm } (f \ x) \bigwedge i. x \in \text{space } M \implies \text{norm } (\text{norm } (s \ (r \ i) \ x)) \leq 2 * \text{norm } (\text{norm } (f \ x))$

using s **by** (*auto intro: LIMSEQ-subseq-LIMSEQ*[*OF tendsto-norm r, unfolded comp-def*] *simple-function-compose1*)

obtain r' **where** r' : *strict-mono* r' **and** *has-cond-exp* $M \ F \ (\lambda x. \text{norm } (f \ x)) \ (\lambda x. \lim (\lambda i. \text{cond-exp } M \ F \ (\lambda x. \text{norm } (s \ (r \ (r' \ i)) \ x)) \ x)) \ AE \ x \text{ in } M. (\lambda i. \text{cond-exp } M \ F \ (\lambda x. \text{norm } (s \ (r \ (r' \ i)) \ x)) \ x) \longrightarrow \lim (\lambda i. \text{cond-exp } M \ F \ (\lambda x. \text{norm } (s \ (r \ (r' \ i)) \ x)) \ x)$ **using** *has-cond-exp-simple-lim*[*OF integrable-norm norm-s-r, OF assms*] **unfolding** *convergent-LIMSEQ-iff* **by** *blast*

hence $r'\text{-tendsto}$: $AE \ x \text{ in } M. (\lambda i. \text{cond-exp } M \ F \ (\lambda x. \text{norm } (s \ (r \ (r' \ i)) \ x)) \ x) \longrightarrow \text{cond-exp } M \ F \ (\lambda x. \text{norm } (f \ x)) \ x$ **using** *has-cond-exp-charact*(2) **by** *force*

have $AE \ x \text{ in } M. \forall i. \text{norm } (\text{cond-exp } M \ F \ (s \ (r \ (r' \ i)) \ x)) \leq \text{cond-exp } M \ F \ (\lambda x. \text{norm } (s \ (r \ (r' \ i)) \ x)) \ x$ **using** s **by** (*auto intro: cond-exp-contraction-simple simp add: AE-all-countable*)

moreover have $AE \ x \text{ in } M. (\lambda i. \text{norm } (\text{cond-exp } M \ F \ (s \ (r \ (r' \ i)) \ x)) \longrightarrow \text{norm } (\text{cond-exp } M \ F \ f \ x))$ **using** $r\text{-tendsto}$ *LIMSEQ-subseq-LIMSEQ*[*OF tendsto-norm r', unfolded comp-def*] **by** *fast*

ultimately show *?thesis* **using** *LIMSEQ-le r'-tendsto* **by** *fast*
qed

lemma *cond-exp-measurable-mult*:

fixes $f \ g :: 'a \Rightarrow \text{real}$

assumes [*measurable*]: *integrable* $M \ (\lambda x. f \ x * g \ x)$ *integrable* $M \ g \ f \in \text{borel-measurable } F$

shows *integrable* $M \ (\lambda x. f \ x * \text{cond-exp } M \ F \ g \ x)$

$AE \ x \text{ in } M. \text{cond-exp } M \ F \ (\lambda x. f \ x * g \ x) \ x = f \ x * \text{cond-exp } M \ F \ g \ x$

proof—

show *integrable*: *integrable* $M \ (\lambda x. f \ x * \text{cond-exp } M \ F \ g \ x)$ **using** *cond-exp-real*[*OF assms*(2)] **by** (*intro integrable-cong-AE-imp*[*OF real-cond-exp-intg*(1), *OF assms*(1,3) *assms*(2)[*THEN borel-measurable-integrable*]] *measurable-from-subalg*[*OF subalg*]) *auto*

interpret *sigma-finite-measure restr-to-subalg* $M \ F$ **by** (*rule sigma-fin-subalg*)


```

{
  fix A assume asm: A ∈ sets F
  hence asm': A ∈ sets M using subalg by (fastforce simp add: subalgebra-def)
  have set-lebesgue-integral M A (cond-exp M F (λx. f x * g x)) = set-lebesgue-integral
M A (λx. f x * g x) by (simp add: cond-exp-set-integral[OF assms(1) asm])
  also have ... = set-lebesgue-integral M A (λx. f x * real-cond-exp M F g
x) using borel-measurable-times[OF borel-measurable-indicator[OF asm] assms(3)]
borel-measurable-integrable[OF assms(2)] integrable-mult-indicator[OF asm' assms(1)]
by (fastforce simp add: set-lebesgue-integral-def mult.assoc[symmetric] intro: real-cond-exp-intg(2)[symmetric])
  also have ... = set-lebesgue-integral M A (λx. f x * cond-exp M F g x) using
cond-exp-real[OF assms(2)] asm' borel-measurable-cond-exp' borel-measurable-cond-exp2
measurable-from-subalg[OF subalg assms(3)] by (auto simp add: set-lebesgue-integral-def
intro: integral-cong-AE)
  finally have set-lebesgue-integral M A (cond-exp M F (λx. f x * g x)) = ∫ x∈A.
(f x * cond-exp M F g x) ∂M .
}
  hence AE x in restr-to-subalg M F. cond-exp M F (λx. f x * g x) x = f
x * cond-exp M F g x by (intro density-unique-banach integrable-cond-exp inte-
grable integrable-in-subalg subalg, measurable, simp add: set-lebesgue-integral-def
integral-subalgebra2[OF subalg] sets-restr-to-subalg[OF subalg])
  thus AE x in M. cond-exp M F (λx. f x * g x) x = f x * cond-exp M F g x by
(rule AE-restr-to-subalg[OF subalg])
qed

```

lemma *cond-exp-measurable-scaleR*:

```

  fixes f :: 'a ⇒ real and g :: 'a ⇒ 'b :: {second-countable-topology, banach}
  assumes [measurable]: integrable M (λx. f x *R g x) integrable M g f ∈ borel-measurable
F
  shows integrable M (λx. f x *R cond-exp M F g x)
    AE x in M. cond-exp M F (λx. f x *R g x) x = f x *R cond-exp M F g x
proof –
  let ?F = restr-to-subalg M F
  have subalg': subalgebra M (restr-to-subalg M F) by (metis sets-eq-imp-space-eq
sets-restr-to-subalg subalg subalgebra-def)
  {
    fix z assume asm[measurable]: integrable M (λx. z x *R g x) z ∈ borel-measurable
?F
    hence asm'[measurable]: z ∈ borel-measurable F using measurable-in-subalg'
subalg by blast
    have integrable M (λx. z x *R cond-exp M F g x) LINT x|M. z x *R g x =
LINT x|M. z x *R cond-exp M F g x
    proof –
    obtain s where s-is: ∧i. simple-function ?F (s i) ∧x. x ∈ space ?F ⇒ (λi.
s i x) ⇒ z x ∧i x. x ∈ space ?F ⇒ norm (s i x) ≤ 2 * norm (z x) using
borel-measurable-implies-sequence-metric[OF asm(2), of 0] by force

```

have *s-scaleR-g-tendsto*: AE x in M. (λi. s i x *_R g x) ⇒ z x *_R g x

using $s\text{-is}(2)$ **by** (*simp add: space-restr-to-subalg tendsto-scaleR*)
have $s\text{-scaleR-cond-exp-g-tendsto}$: $AE\ x\ in\ ?F. (\lambda i. s\ i\ x\ *_R\ cond\text{-}exp\ M\ F\ g\ x) \longrightarrow z\ x\ *_R\ cond\text{-}exp\ M\ F\ g\ x$ **using** $s\text{-is}(2)$ **by** (*simp add: tendsto-scaleR*)

have $s\text{-scaleR-g-meas}$: $(\lambda x. s\ i\ x\ *_R\ g\ x) \in borel\text{-}measurable\ M$ **for** i **using** $s\text{-is}(1)$ [*THEN borel-measurable-simple-function, THEN subalg' [THEN measurable-from-subalg]*] **by** *simp*

have $s\text{-scaleR-cond-exp-g-meas}$: $(\lambda x. s\ i\ x\ *_R\ cond\text{-}exp\ M\ F\ g\ x) \in borel\text{-}measurable\ ?F$ **for** i **using** $s\text{-is}(1)$ [*THEN borel-measurable-simple-function*] *measurable-in-subalg [OF subalg borel-measurable-cond-exp]* **by** (*fastforce intro: borel-measurable-scaleR*)

have $s\text{-scaleR-g-AE-bdd}$: $AE\ x\ in\ M. norm\ (s\ i\ x\ *_R\ g\ x) \leq 2 * norm\ (z\ x\ *_R\ g\ x)$ **for** i **using** $s\text{-is}(3)$ **by** (*fastforce simp add: space-restr-to-subalg mult.assoc[symmetric] mult-right-mono*)

{
fix i
have asm : *integrable* $M\ (\lambda x. norm\ (z\ x) * norm\ (g\ x))$ **using** $asm(1)$ [*THEN integrable-norm*] **by** *simp*
have $AE\ x\ in\ ?F. norm\ (s\ i\ x\ *_R\ cond\text{-}exp\ M\ F\ g\ x) \leq 2 * norm\ (z\ x) * norm\ (cond\text{-}exp\ M\ F\ g\ x)$ **using** $s\text{-is}(3)$ **by** (*fastforce simp add: mult-mono*)
moreover **have** $AE\ x\ in\ ?F. norm\ (z\ x) * cond\text{-}exp\ M\ F\ (\lambda x. norm\ (g\ x))\ x = cond\text{-}exp\ M\ F\ (\lambda x. norm\ (z\ x) * norm\ (g\ x))\ x$ **by** (*rule cond-exp-measurable-mult(2) [THEN AE-symmetric, OF asm integrable-norm, OF assms(2), THEN AE-restr-to-subalg2 [OF subalg]], auto*)
ultimately **have** $AE\ x\ in\ ?F. norm\ (s\ i\ x\ *_R\ cond\text{-}exp\ M\ F\ g\ x) \leq 2 * cond\text{-}exp\ M\ F\ (\lambda x. norm\ (z\ x\ *_R\ g\ x))\ x$ **using** *cond-exp-contraction [OF assms(2), THEN AE-restr-to-subalg2 [OF subalg]] order-trans [OF - mult-mono]* **by** *fastforce*
 }

note $s\text{-scaleR-cond-exp-g-AE-bdd} = this$

{
fix i
have $s\text{-meas-M[measurable]}$: $s\ i \in borel\text{-}measurable\ M$ **by** (*meson borel-measurable-simple-function measurable-from-subalg s-is(1) subalg'*)
have $s\text{-meas-F[measurable]}$: $s\ i \in borel\text{-}measurable\ F$ **by** (*meson borel-measurable-simple-function measurable-in-subalg' s-is(1) subalg*)

have $s\text{-scaleR-eq}$: $s\ i\ x\ *_R\ h\ x = (\sum_{y \in s\ i\ 'space\ M. (indicator\ (s\ i\ -\ ' \{y\} \cap space\ M)\ x\ *_R\ y) *_R\ h\ x})$ **if** $x \in space\ M$ **for** x **and** $h :: 'a \Rightarrow 'b$ **using** *simple-function-indicator-representation [OF s-is(1), of x i]* *that unfolding space-restr-to-subalg scaleR-left.sum [of - h x, symmetric] by presburger*

have $LINT\ x|M. s\ i\ x\ *_R\ g\ x = LINT\ x|M. (\sum_{y \in s\ i\ 'space\ M. indicator\ (s\ i\ -\ ' \{y\} \cap space\ M)\ x\ *_R\ y *_R\ g\ x})$ **using** $s\text{-scaleR-eq}$ **by** (*intro Bochner-Integration.integral-cong*) *auto*

also **have** $... = (\sum_{y \in s\ i\ 'space\ M. LINT\ x|M. indicator\ (s\ i\ -\ ' \{y\} \cap space\ M)\ x\ *_R\ y *_R\ g\ x})$ **by** (*intro Bochner-Integration.integral-sum in-*

tegrable-mult-indicator[*OF* - *integrable-scaleR-right*] *assms*(2)) *simp*
also have ... = ($\sum y \in s \text{ } i \text{ ' } \text{space } M. \text{ } y *_R \text{ set-lebesgue-integral } M \text{ } (s \text{ } i \text{ - ' } \{y\} \cap \text{space } M) \text{ } g$) **by** (*simp only*: *set-lebesgue-integral-def*[*symmetric*]) *simp*
also have ... = ($\sum y \in s \text{ } i \text{ ' } \text{space } M. \text{ } y *_R \text{ set-lebesgue-integral } M \text{ } (s \text{ } i \text{ - ' } \{y\} \cap \text{space } M) \text{ } (\text{cond-exp } M \text{ } F \text{ } g)$) **using** *assms*(2) *subalg borel-measurable-vimage*[*OF s-meas-F*] **by** (*subst cond-exp-set-integral*, *auto simp add*: *subalgebra-def*)
also have ... = ($\sum y \in s \text{ } i \text{ ' } \text{space } M. \text{ } LINT \text{ } x | M. \text{ indicator } (s \text{ } i \text{ - ' } \{y\} \cap \text{space } M) \text{ } x *_R y *_R \text{ cond-exp } M \text{ } F \text{ } g \text{ } x$) **by** (*simp only*: *set-lebesgue-integral-def*[*symmetric*]) *simp*
also have ... = *LINT* *x* | *M*. ($\sum y \in s \text{ } i \text{ ' } \text{space } M. \text{ indicator } (s \text{ } i \text{ - ' } \{y\} \cap \text{space } M) \text{ } x *_R y *_R \text{ cond-exp } M \text{ } F \text{ } g \text{ } x$) **by** (*intro Bochner-Integration.integral-sum*[*symmetric*] *integrable-mult-indicator*[*OF* - *integrable-scaleR-right*]) *auto*
also have ... = *LINT* *x* | *M*. *s i x *_R cond-exp M F g x* **using** *s-scaleR-eq*
by (*intro Bochner-Integration.integral-cong*) *auto*
finally have *LINT* *x* | *M*. *s i x *_R g x* = *LINT* *x* | ?*F*. *s i x *_R cond-exp M F g x* **by** (*simp add*: *integral-subalgebra2*[*OF subalg*])
}
note *integral-s-eq* = *this*

show *integrable* *M* ($\lambda x. z \text{ } x *_R \text{ cond-exp } M \text{ } F \text{ } g \text{ } x$) **using** *s-scaleR-cond-exp-g-meas asm*(2) *borel-measurable-cond-exp'* **by** (*intro integrable-from-subalg*[*OF subalg*] *integrable-cond-exp integrable-dominated-convergence*[*OF* - - - *s-scaleR-cond-exp-g-tendsto s-scaleR-cond-exp-g-AE-bdd*]) (*auto intro*: *measurable-from-subalg*[*OF subalg*] *integrable-in-subalg measurable-in-subalg subalg*)

have ($\lambda i. \text{ } LINT \text{ } x | M. \text{ } s \text{ } i \text{ } x *_R \text{ } g \text{ } x$) \longrightarrow *LINT* *x* | *M*. *z x *_R g x* **using** *s-scaleR-g-meas asm*(1)[*THEN integrable-norm*] *asm'* *borel-measurable-cond-exp'* **by** (*intro integral-dominated-convergence*[*OF* - - - *s-scaleR-g-tendsto s-scaleR-g-AE-bdd*]) (*auto intro*: *measurable-from-subalg*[*OF subalg*])

moreover have ($\lambda i. \text{ } LINT \text{ } x | ?F. \text{ } s \text{ } i \text{ } x *_R \text{ cond-exp } M \text{ } F \text{ } g \text{ } x$) \longrightarrow *LINT* *x* | ?*F*. *z x *_R cond-exp M F g x* **using** *s-scaleR-cond-exp-g-meas asm*(2) *borel-measurable-cond-exp'* **by** (*intro integral-dominated-convergence*[*OF* - - - *s-scaleR-cond-exp-g-tendsto s-scaleR-cond-exp-g-AE-bdd*]) (*auto intro*: *measurable-from-subalg*[*OF subalg*] *integrable-in-subalg measurable-in-subalg subalg*)

ultimately show *LINT* *x* | *M*. *z x *_R g x* = *LINT* *x* | *M*. *z x *_R cond-exp M F g x* **using** *integral-s-eq using subalg* **by** (*simp add*: *LIMSEQ-unique integral-subalgebra2*)

qed
}
note * = *this*

show *integrable* *M* ($\lambda x. f \text{ } x *_R \text{ cond-exp } M \text{ } F \text{ } g \text{ } x$) **using** * *assms measurable-in-subalg*[*OF subalg*] **by** *blast*

{

fix A **assume** $asm: A \in F$
hence $integrable\ M\ (\lambda x. indicat-real\ A\ x\ *_R\ f\ x\ *_R\ g\ x)$ **using** $subalg$ **by**
(fastforce simp add: subalgebra-def intro!: integrable-mult-indicator assms(1))
hence $set-lebesgue-integral\ M\ A\ (\lambda x. f\ x\ *_R\ g\ x) = set-lebesgue-integral\ M\ A$
 $(\lambda x. f\ x\ *_R\ cond-exp\ M\ F\ g\ x)$ **unfolding** $set-lebesgue-integral-def$ **using** asm **by**
*(auto intro!: * measurable-in-subalg[OF subalg])*
}
thus $AE\ x\ in\ M. cond-exp\ M\ F\ (\lambda x. f\ x\ *_R\ g\ x)\ x = f\ x\ *_R\ cond-exp\ M\ F\ g\ x$
using $borel-measurable-cond-exp$ **by** *(intro cond-exp-charact, auto intro!: * assms measurable-in-subalg[OF subalg])*
qed

lemma $cond-exp-sum$ *[intro, simp]:*
fixes $f :: 't \Rightarrow 'a \Rightarrow 'b :: \{second-countable-topology, banach\}$
assumes *[measurable]:* $\bigwedge i. integrable\ M\ (f\ i)$
shows $AE\ x\ in\ M. cond-exp\ M\ F\ (\lambda x. \sum_{i \in I}. f\ i\ x)\ x = (\sum_{i \in I}. cond-exp\ M\ F\ (f\ i)\ x)$
proof *(rule has-cond-exp-charact, intro has-cond-expI')*
fix A **assume** *[measurable]:* $A \in sets\ F$
then have $A-meas\ [measurable]: A \in sets\ M$ **by** *(meson subsetD subalg subalgebra-def)*

have $(\int x \in A. (\sum_{i \in I}. f\ i\ x) \partial M) = (\int x. (\sum_{i \in I}. indicator\ A\ x\ *_R\ f\ i\ x) \partial M)$
unfolding $set-lebesgue-integral-def$ **by** *(simp add: scaleR-sum-right)*
also have $\dots = (\sum_{i \in I}. (\int x. indicator\ A\ x\ *_R\ f\ i\ x\ \partial M))$ **using** $assms$ **by** *(auto intro!: Bochner-Integration.integral-sum integrable-mult-indicator)*
also have $\dots = (\sum_{i \in I}. (\int x. indicator\ A\ x\ *_R\ cond-exp\ M\ F\ (f\ i)\ x\ \partial M))$ **using**
 $cond-exp-set-integral[OF assms]$ **by** *(simp add: set-lebesgue-integral-def)*
also have $\dots = (\int x. (\sum_{i \in I}. indicator\ A\ x\ *_R\ cond-exp\ M\ F\ (f\ i)\ x) \partial M)$
using $assms$ **by** *(auto intro!: Bochner-Integration.integral-sum[symmetric] integrable-mult-indicator)*
also have $\dots = (\int x \in A. (\sum_{i \in I}. cond-exp\ M\ F\ (f\ i)\ x) \partial M)$ **unfolding** $set-lebesgue-integral-def$
by *(simp add: scaleR-sum-right)*
finally show $(\int x \in A. (\sum_{i \in I}. f\ i\ x) \partial M) = (\int x \in A. (\sum_{i \in I}. cond-exp\ M\ F\ (f\ i)\ x) \partial M)$ **by** $auto$
qed *(auto simp add: assms integrable-cond-exp)*

4.3 Linearly Ordered Banach Spaces

In this subsection we show monotonicity results concerning the conditional expectation operator.

lemma $cond-exp-gr-c$:
fixes $f :: 'a \Rightarrow 'b :: \{second-countable-topology, banach, linorder-topology, ordered-real-vector\}$
assumes $integrable\ M\ f\ AE\ x\ in\ M. f\ x > c$
shows $AE\ x\ in\ M. cond-exp\ M\ F\ f\ x > c$
proof –
define X **where** $X = \{x \in space\ M. cond-exp\ M\ F\ f\ x \leq c\}$
have *[measurable]:* $X \in sets\ F$ **unfolding** $X-def$ **by** $measurable\ (metis\ sets.top)$

subalg subalgebra-def)
hence $X\text{-in-}M$: $X \in \text{sets } M$ **using** *sets-restr-to-subalg subalg subalgebra-def* **by**
blast
have $\text{emeasure } M X = 0$
proof (*rule ccontr*)
assume $\text{emeasure } M X \neq 0$
have $\text{emeasure } (\text{restr-to-subalg } M F) X = \text{emeasure } M X$ **by** (*simp add: emea-*
sure-restr-to-subalg subalg)
hence $\text{emeasure } (\text{restr-to-subalg } M F) X > 0$ **using** $\neg(\text{emeasure } M X) = 0$
gr-zeroI **by** *auto*
then obtain A **where** $A: A \in \text{sets } (\text{restr-to-subalg } M F) A \subseteq X$ emeasure
 $(\text{restr-to-subalg } M F) A > 0$ $\text{emeasure } (\text{restr-to-subalg } M F) A < \infty$
using *sigma-fin-subalg* **by** (*metis emeasure-notin-sets ennreal-0 infinity-ennreal-def*
le-less-linear neq-top-trans not-gr-zero order-refl sigma-finite-measure.approx-PInf-emeasure-with-finite)
hence $[simp]: A \in \text{sets } F$ **using** *subalg sets-restr-to-subalg* **by** *blast*
hence $A\text{-in-sets-}M[simp]: A \in \text{sets } M$ **using** *sets-restr-to-subalg subalg subal-*
gebra-def **by** *blast*
have $[simp]: \text{set-integrable } M A (\lambda x. c)$ **using** A *subalg* **by** (*auto simp add:*
set-integrable-def emeasure-restr-to-subalg)
have $[simp]: \text{set-integrable } M A f$ **unfolding** *set-integrable-def* **by** (*rule inte-*
grable-mult-indicator, auto simp add: assms(1))
have $AE\ x\ \text{in}\ M. \text{indicator } A\ x *_R c = \text{indicator } A\ x *_R f\ x$
proof (*rule integral-eq-mono-AE-eq-AE*)
have $(\int x \in A. c\ \partial M) \leq (\int x \in A. f\ x\ \partial M)$ **using** *assms(2)* **by** (*intro set-integral-mono-AE-banach*)
auto
moreover
{
have $(\int x \in A. f\ x\ \partial M) = (\int x \in A. \text{cond-exp } M\ F\ f\ x\ \partial M)$ **by** (*rule*
cond-exp-set-integral, auto simp add: assms)
also have $\dots \leq (\int x \in A. c\ \partial M)$ **using** A **by** (*auto intro!: set-integral-mono-banach*
simp add: X-def)
finally have $(\int x \in A. f\ x\ \partial M) \leq (\int x \in A. c\ \partial M)$ **by** *simp*
}
ultimately show $LINT\ x|M. \text{indicator } A\ x *_R c = LINT\ x|M. \text{indicator } A$
 $x *_R f\ x$ **unfolding** *set-lebesgue-integral-def* **by** *simp*
show $AE\ x \in A\ \text{in}\ M. \text{indicator } A\ x *_R c \leq \text{indicator } A\ x *_R f\ x$ **using** *assms* **by**
(auto simp add: X-def indicator-def)
qed (*auto simp add: set-integrable-def[symmetric]*)
hence $AE\ x \in A\ \text{in}\ M. c = f\ x$ **by** *auto*
hence $AE\ x \in A\ \text{in}\ M. \text{False}$ **using** *assms(2)* **by** *auto*
hence $A \in \text{null-sets } M$ **using** *AE-iff-null-sets A-in-sets-M* **by** *metis*
thus False **using** *A(3)* **by** (*simp add: emeasure-restr-to-subalg null-setsD1*
subalg)
qed
thus *?thesis* **using** *AE-iff-null-sets[OF X-in-M]* **unfolding** *X-def* **by** *auto*
qed

corollary *cond-exp-less-c*:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$

```

dered-real-vector}
  assumes integrable M f AE x in M. f x < c
  shows AE x in M. cond-exp M F f x < c
proof -
  have AE x in M. cond-exp M F f x = - cond-exp M F (λx. - f x) x using
cond-exp-uminus[OF assms(1)] by auto
  moreover have AE x in M. cond-exp M F (λx. - f x) x > - c using assms
by (intro cond-exp-gr-c) auto
  ultimately show ?thesis by (force simp add: minus-less-iff)
qed

```

lemma *cond-exp-mono-strict*:

```

  fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
dered-real-vector}
  assumes integrable M f integrable M g AE x in M. f x < g x
  shows AE x in M. cond-exp M F f x < cond-exp M F g x
  using cond-exp-less-c[OF Bochner-Integration.integrable-diff, OF assms(1,2), of
0]
cond-exp-diff[OF assms(1,2)] assms(3) by auto

```

lemma *cond-exp-ge-c*:

```

  fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
dered-real-vector}
  assumes [measurable]: integrable M f
  and AE x in M. f x ≥ c
  shows AE x in M. cond-exp M F f x ≥ c
proof -
  let ?F = restr-to-subalg M F
  interpret sigma-finite-measure restr-to-subalg M F using sigma-fin-subalg by
auto
  {
    fix A assume asm: A ∈ sets ?F 0 < measure ?F A
    have [simp]: sets ?F = sets F measure ?F A = measure M A using asm by (auto
simp add: measure-def sets-restr-to-subalg[OF subalg] emeasure-restr-to-subalg[OF
subalg])
    have M-A: emeasure M A < ∞ using measure-zero-top asm by (force simp
add: top.not-eq-extremum)
    hence F-A: emeasure ?F A < ∞ using asm(1) emeasure-restr-to-subalg subalg
by fastforce
    have set-lebesgue-integral M A (λ-. c) ≤ set-lebesgue-integral M A f using
assms asm M-A subalg by (intro set-integral-mono-AE-banach, auto simp add:
set-integrable-def integrable-mult-indicator subalgebra-def sets-restr-to-subalg)
    also have ... = set-lebesgue-integral M A (cond-exp M F f) using cond-exp-set-integral[OF
assms(1)] asm by auto
    also have ... = set-lebesgue-integral ?F A (cond-exp M F f) unfolding set-lebesgue-integral-def
using asm borel-measurable-cond-exp by (intro integral-subalgebra2[OF subalg, sym-
metric], simp)
    finally have (1 / measure ?F A) *R set-lebesgue-integral ?F A (cond-exp M F f)
∈ {c..} using asm subalg M-A by (auto simp add: set-integral-const subalgebra-def

```

intro!: pos-divideR-le-eq[THEN iffD1])
 }
 thus ?thesis using AE-restr-to-subalg[OF subalg] averaging-theorem[OF inte-
 grable-in-subalg closed-atLeast, OF subalg borel-measurable-cond-exp integrable-cond-exp]
 by auto
 qed

corollary cond-exp-le-c:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$
 $\text{dered-real-vector}\}$
 assumes integrable $M f$
 and $AE\ x\ in\ M. f\ x \leq c$
 shows $AE\ x\ in\ M. \text{cond-exp } M F f\ x \leq c$
proof –
 have $AE\ x\ in\ M. \text{cond-exp } M F f\ x = - \text{cond-exp } M F (\lambda x. - f\ x)\ x$ using
 cond-exp-uminus[OF assms(1)] by force
 moreover have $AE\ x\ in\ M. \text{cond-exp } M F (\lambda x. - f\ x)\ x \geq - c$ using assms
 by (intro cond-exp-ge-c) auto
 ultimately show ?thesis by (force simp add: minus-le-iff)
 qed

corollary cond-exp-mono:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$
 $\text{dered-real-vector}\}$
 assumes integrable $M f$ integrable $M g$ $AE\ x\ in\ M. f\ x \leq g\ x$
 shows $AE\ x\ in\ M. \text{cond-exp } M F f\ x \leq \text{cond-exp } M F g\ x$
 using cond-exp-le-c[OF Bochner-Integration.integrable-diff, OF assms(1,2), of
 0]
 cond-exp-diff[OF assms(1,2)] assms(3) by auto

corollary cond-exp-min:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$
 $\text{dered-real-vector}\}$
 assumes integrable $M f$ integrable $M g$
 shows $AE\ \xi\ in\ M. \text{cond-exp } M F (\lambda x. \min (f\ x) (g\ x))\ \xi \leq \min (\text{cond-exp } M F$
 $f\ \xi) (\text{cond-exp } M F g\ \xi)$
proof –
 have $AE\ \xi\ in\ M. \text{cond-exp } M F (\lambda x. \min (f\ x) (g\ x))\ \xi \leq \text{cond-exp } M F f\ \xi$ by
 (intro cond-exp-mono integrable-min assms, simp)
 moreover have $AE\ \xi\ in\ M. \text{cond-exp } M F (\lambda x. \min (f\ x) (g\ x))\ \xi \leq \text{cond-exp}$
 $M F g\ \xi$ by (intro cond-exp-mono integrable-min assms, simp)
 ultimately show $AE\ \xi\ in\ M. \text{cond-exp } M F (\lambda x. \min (f\ x) (g\ x))\ \xi \leq \min$
 $(\text{cond-exp } M F f\ \xi) (\text{cond-exp } M F g\ \xi)$ by fastforce
 qed

corollary cond-exp-max:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, linorder-topology, or-}$
 $\text{dered-real-vector}\}$
 assumes integrable $M f$ integrable $M g$

shows $AE \xi$ in M . $cond\text{-}exp\ M\ F\ (\lambda x. \max (f\ x)\ (g\ x))\ \xi \geq \max (cond\text{-}exp\ M\ F\ f\ \xi)\ (cond\text{-}exp\ M\ F\ g\ \xi)$

proof –

have $AE \xi$ in M . $cond\text{-}exp\ M\ F\ (\lambda x. \max (f\ x)\ (g\ x))\ \xi \geq cond\text{-}exp\ M\ F\ f\ \xi$ **by** $(intro\ cond\text{-}exp\text{-}mono\ integrable\text{-}max\ assms, simp)$

moreover have $AE \xi$ in M . $cond\text{-}exp\ M\ F\ (\lambda x. \max (f\ x)\ (g\ x))\ \xi \geq cond\text{-}exp\ M\ F\ g\ \xi$ **by** $(intro\ cond\text{-}exp\text{-}mono\ integrable\text{-}max\ assms, simp)$

ultimately show $AE \xi$ in M . $cond\text{-}exp\ M\ F\ (\lambda x. \max (f\ x)\ (g\ x))\ \xi \geq \max (cond\text{-}exp\ M\ F\ f\ \xi)\ (cond\text{-}exp\ M\ F\ g\ \xi)$ **by** $fastforce$

qed

corollary $cond\text{-}exp\text{-}inf$:

fixes $f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, ordered\text{-}real\text{-}vector, lattice\}$

assumes $integrable\ M\ f\ integrable\ M\ g$

shows $AE \xi$ in M . $cond\text{-}exp\ M\ F\ (\lambda x. \inf (f\ x)\ (g\ x))\ \xi \leq \inf (cond\text{-}exp\ M\ F\ f\ \xi)\ (cond\text{-}exp\ M\ F\ g\ \xi)$

unfolding $inf\text{-}min$ **using** $assms$ **by** $(rule\ cond\text{-}exp\text{-}min)$

corollary $cond\text{-}exp\text{-}sup$:

fixes $f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach, linorder\text{-}topology, ordered\text{-}real\text{-}vector, lattice\}$

assumes $integrable\ M\ f\ integrable\ M\ g$

shows $AE \xi$ in M . $cond\text{-}exp\ M\ F\ (\lambda x. \sup (f\ x)\ (g\ x))\ \xi \geq \sup (cond\text{-}exp\ M\ F\ f\ \xi)\ (cond\text{-}exp\ M\ F\ g\ \xi)$

unfolding $sup\text{-}max$ **using** $assms$ **by** $(rule\ cond\text{-}exp\text{-}max)$

end

4.4 Probability Spaces

lemma (in $prob\text{-}space$) $sigma\text{-finite}\text{-}subalgebra\text{-}restr\text{-}to\text{-}subalg$:

assumes $subalgebra\ M\ F$

shows $sigma\text{-finite}\text{-}subalgebra\ M\ F$

proof $(intro\ sigma\text{-finite}\text{-}subalgebra.intro)$

interpret F : $prob\text{-}space\ restr\text{-}to\text{-}subalg\ M\ F$ **using** $assms\ prob\text{-}space\text{-}restr\text{-}to\text{-}subalg$

$prob\text{-}space\text{-}axioms$ **by** $blast$

show $sigma\text{-finite}\text{-}measure\ (restr\text{-}to\text{-}subalg\ M\ F)$ **by** $(rule\ F.sigma\text{-finite}\text{-}measure\text{-}axioms)$

qed $(rule\ assms)$

lemma (in $prob\text{-}space$) $cond\text{-}exp\text{-}trivial$:

fixes $f :: 'a \Rightarrow 'b :: \{second\text{-}countable\text{-}topology, banach\}$

assumes $integrable\ M\ f$

shows $AE\ x$ in M . $cond\text{-}exp\ M\ (sigma\ (space\ M)\ \{\})\ f\ x = expectation\ f$

proof –

interpret $sigma\text{-finite}\text{-}subalgebra\ M\ sigma\ (space\ M)\ \{\}$ **by** $(auto\ intro: sigma\text{-finite}\text{-}subalgebra\text{-}restr\text{-}to\text{-}subalg\ simp\ add: subalgebra\text{-}def\ sigma\text{-sets}\text{-}empty\text{-}eq)$

show $?thesis$ **using** $assms$ **by** $(intro\ cond\text{-}exp\text{-}charact)\ (auto\ simp\ add: sigma\text{-sets}\text{-}empty\text{-}eq\ set\text{-}lebesgue\text{-}integral\text{-}def\ prob\text{-}space\ cong: Bochner\text{-}Integration.integral\text{-}cong)$

qed

The following lemma shows that independent sigma algebras don't matter for the conditional expectation.

lemma (in prob-space) cond-exp-indep-subalgebra:

fixes $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology}, \text{banach}, \text{real-normed-field}\}$

assumes subalgebra: subalgebra M F subalgebra M G

and independent: indep-set G (sigma (space M) ($F \cup \text{vimage-algebra (space } M) f \text{ borel}$))

assumes [measurable]: integrable M f

shows AE x in M . cond-exp M (sigma (space M) ($F \cup G$)) f $x = \text{cond-exp } M$ F f x

proof –

interpret Un-sigma: sigma-finite-subalgebra M sigma (space M) ($F \cup G$) **using** $\text{assms}(1,2)$ **by** (auto intro!: sigma-finite-subalgebra-restr-to-subalg sets.sigma-sets-subset simp add: subalgebra-def space-measure-of-conv sets-measure-of-conv)

interpret sigma-finite-subalgebra M F **using** assms **by** (auto intro: sigma-finite-subalgebra-restr-to-subalg)

{

fix A

assume $\text{asm}: A \in \text{sigma (space } M) \{a \cap b \mid a \in F \wedge b \in G\}$

have in-events: sigma-sets (space M) $\{a \cap b \mid a \in \text{sets } F \wedge b \in \text{sets } G\} \subseteq \text{events}$ **using** subalgebra **by** (intro sets.sigma-sets-subset, auto simp add: subalgebra-def)

have Int-stable $\{a \cap b \mid a \in F \wedge b \in G\}$

proof –

{

fix af bf ag bg

assume $F: af \in F$ $bf \in F$ **and** $G: ag \in G$ $bg \in G$

have $af \cap bf \in F$ **by** (intro sets.Int F)

moreover have $ag \cap bg \in G$ **by** (intro sets.Int G)

ultimately have $\exists a \in F. af \cap ag \cap (bf \cap bg) = a \cap b \wedge a \in \text{sets } F \wedge b \in \text{sets } G$ **by** (metis inf-assoc inf-left-commute)

}

thus ?thesis **by** (force intro!: Int-stableI)

qed

moreover have $\{a \cap b \mid a \in F \wedge b \in G\} \subseteq \text{Pow (space } M)$ **using** subalgebra **by** (force simp add: subalgebra-def dest: sets.sets-into-space)

moreover have $A \in \text{sigma-sets (space } M) \{a \cap b \mid a \in F \wedge b \in G\}$ **using** calculation asm **by** force

ultimately have set-lebesgue-integral M A $f = \text{set-lebesgue-integral } M$ A (cond-exp M F f)

proof (induction rule: sigma-sets-induct-disjoint)

case (basic A)

then obtain $a \in F$ $b \in G$ **where** $A = a \cap b$ **by** blast

hence events[measurable]: $a \in \text{events}$ $b \in \text{events}$ **using** subalgebra **by** (auto simp add: subalgebra-def)

have [simp]: sigma-sets (space M) $\{\text{indicator } b - 'A \cap \text{space } M \mid A \in \text{borel}\}$

$\subseteq G$
using *borel-measurable-indicator*[*OF* *A*(3), *THEN* *measurable-sets*] *sets.top*
subalgebra
by (*intro sets.sigma-sets-subset'*) (*fastforce simp add: subalgebra-def*)+

have *Un-in-sigma*: $F \cup \text{vimage-algebra } (\text{space } M) f \text{ borel} \subseteq \text{sigma } (\text{space } M) (F \cup \text{vimage-algebra } (\text{space } M) f \text{ borel})$ **by** (*metis equalityE le-supI sets.space-closed sigma-le-sets space-vimage-algebra subalg subalgebra-def*)

have [*intro*]: *indep-var borel (indicator b) borel* ($\lambda \omega. \text{indicator } a \ \omega *_{\mathcal{R}} f \ \omega$)
proof –
have [*simp*]: $\text{sigma-sets } (\text{space } M) \{(\lambda \omega. \text{indicator } a \ \omega *_{\mathcal{R}} f \ \omega) - 'A \cap \text{space } M \mid A. A \in \text{borel}\} \subseteq \text{sigma } (\text{space } M) (F \cup \text{vimage-algebra } (\text{space } M) f \text{ borel})$
proof –
have *: $(\lambda \omega. \text{indicator } a \ \omega *_{\mathcal{R}} f \ \omega) \in \text{borel-measurable } (\text{sigma } (\text{space } M) (F \cup \text{vimage-algebra } (\text{space } M) f \text{ borel}))$
using *borel-measurable-indicator*[*OF* *A*(2), *THEN* *measurable-sets*, *OF* *borel-open*] *subalgebra*
by (*intro borel-measurable-scaleR borel-measurableI Un-in-sigma[THEN subsetD]*)
(*auto simp add: space-measure-of-conv subalgebra-def sets-vimage-algebra2*)
thus ?thesis **using** *measurable-sets*[*OF* *] **by** (*intro sets.sigma-sets-subset'*, *auto simp add: space-measure-of-conv*)
qed
have *indep-set* ($\text{sigma-sets } (\text{space } M) \{\text{indicator } b - 'A \cap \text{space } M \mid A. A \in \text{borel}\}$) ($\text{sigma-sets } (\text{space } M) \{(\lambda \omega. \text{indicator } a \ \omega *_{\mathcal{R}} f \ \omega) - 'A \cap \text{space } M \mid A. A \in \text{borel}\}$)
using *independent unfolding indep-set-def* **by** (*rule indep-sets-mono-sets*, *auto split: bool.split*)
thus ?thesis **by** (*subst indep-var-eq*, *auto intro!: borel-measurable-scaleR*)
qed

have [*intro*]: *indep-var borel (indicator b) borel* ($\lambda \omega. \text{indicat-real } a \ \omega *_{\mathcal{R}}$
cond-exp M F f ω)
proof –
have [*simp*]: $\text{sigma-sets } (\text{space } M) \{(\lambda \omega. \text{indicator } a \ \omega *_{\mathcal{R}} \text{cond-exp } M F f \ \omega) - 'A \cap \text{space } M \mid A. A \in \text{borel}\} \subseteq \text{sigma } (\text{space } M) (F \cup \text{vimage-algebra } (\text{space } M) f \text{ borel})$
proof –
have *: $(\lambda \omega. \text{indicator } a \ \omega *_{\mathcal{R}} \text{cond-exp } M F f \ \omega) \in \text{borel-measurable } (\text{sigma } (\text{space } M) (F \cup \text{vimage-algebra } (\text{space } M) f \text{ borel}))$
using *borel-measurable-indicator*[*OF* *A*(2), *THEN* *measurable-sets*, *OF* *borel-open*] *subalgebra*
borel-measurable-cond-exp[*THEN* *measurable-sets*, *OF* *borel-open*, *of* *- M F f*]
by (*intro borel-measurable-scaleR borel-measurableI Un-in-sigma[THEN subsetD]*)
(*auto simp add: space-measure-of-conv subalgebra-def*)
thus ?thesis **using** *measurable-sets*[*OF* *] **by** (*intro sets.sigma-sets-subset'*,

```

auto simp add: space-measure-of-conv)
qed
have indep-set (sigma-sets (space M) {indicator b - ' A ∩ space M | A. A ∈
borel}) (sigma-sets (space M) {(λω. indicator a ω *R cond-exp M F f ω) - ' A ∩
space M | A. A ∈ borel})
using independent unfolding indep-set-def by (rule indep-sets-mono-sets,
auto split: bool.split)
thus ?thesis by (subst indep-var-eq, auto intro!: borel-measurable-scaleR)
qed

have set-lebesgue-integral M A f = (LINT x|M. indicator b x * (indicator a
x *R f x))
unfolding set-lebesgue-integral-def A indicator-inter-arith
by (intro Bochner-Integration.integral-cong, auto simp add: scaleR-scaleR[symmetric]
indicator-times-eq-if(1))
also have ... = (LINT x|M. indicator b x) * (LINT x|M. indicator a x *R f
x)
by (intro indep-var-lebesgue-integral
Bochner-Integration.integrable-bound[OF integrable-const[of 1 :: 'b]
borel-measurable-indicator]
integrable-mult-indicator[OF - assms(4)], blast) (auto simp add:
indicator-def)
also have ... = (LINT x|M. indicator b x) * (LINT x|M. indicator a x *R
cond-exp M F f x)
using cond-exp-set-integral[OF assms(4) A(2)] unfolding set-lebesgue-integral-def
by argo
also have ... = (LINT x|M. indicator b x * (indicator a x *R cond-exp M
F f x))
by (intro indep-var-lebesgue-integral[symmetric]
Bochner-Integration.integrable-bound[OF integrable-const[of 1 :: 'b]
borel-measurable-indicator]
integrable-mult-indicator[OF - integrable-cond-exp], blast) (auto simp
add: indicator-def)
also have ... = set-lebesgue-integral M A (cond-exp M F f)
unfolding set-lebesgue-integral-def A indicator-inter-arith
by (intro Bochner-Integration.integral-cong, auto simp add: scaleR-scaleR[symmetric]
indicator-times-eq-if(1))
finally show ?case .
next
case empty
then show ?case unfolding set-lebesgue-integral-def by simp
next
case (compl A)
have A-in-space: A ⊆ space M using compl using in-events sets.sets-into-space
by blast
have set-lebesgue-integral M (space M - A) f = set-lebesgue-integral M (space
M - A ∪ A) f - set-lebesgue-integral M A f
using compl(1) in-events
by (subst set-integral-Un[of space M - A A], blast)

```

```

      (simp | intro integrable-mult-indicator[folded set-integrable-def, OF -
assms(4)], fast)+
    also have ... = set-lebesgue-integral M (space M - A ∪ A) (cond-exp M F f)
  - set-lebesgue-integral M A (cond-exp M F f)
    using cond-exp-set-integral[OF assms(4) sets.top] compl subalgebra by (simp
add: subalgebra-def Un-absorb2[OF A-in-space])
    also have ... = set-lebesgue-integral M (space M - A) (cond-exp M F f)
    using compl(1) in-events
    by (subst set-integral-Un[of space M - A A], blast)
      (simp | intro integrable-mult-indicator[folded set-integrable-def, OF -
integrable-cond-exp], fast)+
    finally show ?case .
  next
    case (union A)
    have set-lebesgue-integral M (⋃ (range A)) f = (∑ i. set-lebesgue-integral M
(A i) f)
      using union in-events
    by (intro lebesgue-integral-countable-add) (auto simp add: disjoint-family-onD
intro!: integrable-mult-indicator[folded set-integrable-def, OF - assms(4)])
    also have ... = (∑ i. set-lebesgue-integral M (A i) (cond-exp M F f)) using
union by presburger
    also have ... = set-lebesgue-integral M (⋃ (range A)) (cond-exp M F f)
    using union in-events
    by (intro lebesgue-integral-countable-add[symmetric]) (auto simp add: dis-
joint-family-onD intro!: integrable-mult-indicator[folded set-integrable-def, OF - in-
tegrable-cond-exp])
    finally show ?case .
  qed
}
moreover have sigma (space M) {a ∩ b | a b. a ∈ F ∧ b ∈ G} = sigma (space
M) (F ∪ G)
proof -
  have sigma-sets (space M) {a ∩ b | a b. a ∈ sets F ∧ b ∈ sets G} = sigma-sets
(space M) (sets F ∪ sets G)
proof -
  {
    fix a b assume asm: a ∈ F b ∈ G
    hence a ∩ b ∈ sigma-sets (space M) (F ∪ G) using subalgebra unfolding
Int-range-binary by (intro sigma-sets-Inter[OF - binary-in-sigma-sets]) (force simp
add: subalgebra-def dest: sets.sets-into-space)+
  }
moreover
  {
    fix a
    assume a ∈ sets F
    hence a ∈ sigma-sets (space M) {a ∩ b | a b. a ∈ sets F ∧ b ∈ sets G}
      using subalgebra sets.top[of G] sets.sets-into-space[of - F]
    by (intro sigma-sets.Basic, auto simp add: subalgebra-def)
  }
}

```

```

moreover
{
  fix  $a$  assume  $a \in \text{sets } F \vee a \in \text{sets } G \wedge a \notin \text{sets } F$ 
  hence  $a \in \text{sets } G$  by blast
  hence  $a \in \text{sigma-sets } (\text{space } M) \{a \cap b \mid a \in \text{sets } F \wedge b \in \text{sets } G\}$ 
    using subalgebra sets.top[of F] sets.sets-into-space[of - G]
    by (intro sigma-sets.Basic, auto simp add: subalgebra-def)
}
ultimately show ?thesis by (intro sigma-sets-eqI) auto
qed
thus ?thesis using subalgebra by (intro sigma-eqI) (force simp add: subalgebra-def dest: sets.sets-into-space)
qed
moreover have (cond-exp M F f)  $\in$  borel-measurable (sigma (space M) (sets F  $\cup$  sets G))
proof -
  have  $F \subseteq \text{sigma } (\text{space } M) (F \cup G)$  by (metis Un-least Un-upper1 measure-of-of-measure sets.space-closed sets-measure-of sigma-sets-subseteq subalg subalgebra(2) subalgebra-def)
  thus ?thesis using borel-measurable-cond-exp[THEN measurable-sets, OF borel-open, of - M F f] subalgebra by (intro borel-measurableI, force simp only: space-measure-of-conv subalgebra-def)
qed
ultimately show ?thesis using assms(4) integrable-cond-exp by (intro Un-sigma.cond-exp-charact presburger)
qed

lemma (in prob-space) cond-exp-indep:
  fixes  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach, real-normed-field}\}$ 
  assumes subalgebra: subalgebra M F
    and independent: indep-set F (vimage-algebra (space M) f borel)
    and integrable: integrable M f
  shows AE x in M. cond-exp M F f x = expectation f
proof -
  have indep-set F (sigma (space M) (sigma (space M) {})  $\cup$  (vimage-algebra (space M) f borel)))
    using independent unfolding indep-set-def
    by (rule indep-sets-mono-sets, simp add: bool.split)
    (metis bot.extremum dual-order.refl sets.sets-measure-of-eq sets.sigma-sets-subset' sets-vimage-algebra-space space-vimage-algebra sup.absorb-iff2)
  hence cond-exp-indep: AE x in M. cond-exp M (sigma (space M) (sigma (space M) {}  $\cup$  F)) f x = expectation f
    using cond-exp-indep-subalgebra[OF - subalgebra - integrable, of sigma (space M) {}] cond-exp-trivial[OF integrable]
    by (auto simp add: subalgebra-def sigma-sets-empty-eq)
  have sets (sigma (space M) (sigma (space M) {}  $\cup$  F)) = F
    using subalgebra sets.top[of F] unfolding subalgebra-def
    by (simp add: sigma-sets-empty-eq, subst insert-absorb[of space M F], blast)
    (metis insert-absorb[OF sets.empty-sets] sets.sets-measure-of-eq)

```

hence $AE\ x\ in\ M.\ cond\text{-}exp\ M\ (\sigma\ (space\ M)\ (\sigma\ (space\ M)\ \{\}\ \cup\ F))\ f$
 $x = cond\text{-}exp\ M\ F\ f\ x$ **by** (*rule cond-exp-sets-cong*)
thus *?thesis* **using** *cond-exp-indep* **by** *force*
qed
end

theory *Filtered-Measure*
imports *HOL-Probability.Conditional-Expectation*
begin

5 Filtered Measure Spaces

5.1 Filtered Measure

locale *filtered-measure* =
fixes $M\ F$ **and** $t_0 :: 'b :: \{second\text{-}countable\text{-}topology,\ order\text{-}topology,\ t2\text{-}space\}$
assumes *subalgebras*: $\bigwedge i.\ t_0 \leq i \implies subalgebra\ M\ (F\ i)$
and *sets-F-mono*: $\bigwedge i\ j.\ t_0 \leq i \implies i \leq j \implies sets\ (F\ i) \leq sets\ (F\ j)$
begin

lemma *space-F[simp]*:
assumes $t_0 \leq i$
shows $space\ (F\ i) = space\ M$
using *subalgebras* *assms* **by** (*simp add: subalgebra-def*)

lemma *subalgebra-F[intro]*:
assumes $t_0 \leq i\ i \leq j$
shows $subalgebra\ (F\ j)\ (F\ i)$
unfolding *subalgebra-def* **using** *assms* **by** (*simp add: sets-F-mono*)

lemma *borel-measurable-mono*:
assumes $t_0 \leq i\ i \leq j$
shows $borel\text{-}measurable\ (F\ i) \subseteq borel\text{-}measurable\ (F\ j)$
unfolding *subset-iff* **by** (*metis assms subalgebra-F measurable-from-subalg*)

end

locale *linearly-filtered-measure* = *filtered-measure* $M\ F\ t_0$ **for** M **and** $F :: - :: \{linorder\text{-}topology\} \Rightarrow -$ **and** t_0

locale *nat-filtered-measure* = *linearly-filtered-measure* $M\ F\ 0$ **for** M **and** $F :: nat \Rightarrow -$

locale *real-filtered-measure* = *linearly-filtered-measure* $M\ F\ 0$ **for** M **and** $F :: real \Rightarrow -$

5.2 Sigma Finite Filtered Measure

The locale presented here is a generalization of the *sigma-finite-subalgebra* for a particular filtration.

locale *sigma-finite-filtered-measure* = *filtered-measure* +
assumes *sigma-finite-initial*: *sigma-finite-subalgebra* *M* (*F* *t*₀)

lemma (**in** *sigma-finite-filtered-measure*) *sigma-finite-subalgebra-F*[*intro*]:
assumes *t*₀ ≤ *i*
shows *sigma-finite-subalgebra* *M* (*F* *i*)
using *assms* **by** (*metis dual-order.refl sets-F-mono sigma-finite-initial sigma-finite-subalgebra.nested-subalgebras subalgebra-def*)

locale *nat-sigma-finite-filtered-measure* = *sigma-finite-filtered-measure* *M* *F* 0 ::
nat **for** *M* *F*

locale *real-sigma-finite-filtered-measure* = *sigma-finite-filtered-measure* *M* *F* 0 ::
real **for** *M* *F*

sublocale *nat-sigma-finite-filtered-measure* ⊆ *sigma-finite-subalgebra* *M* *F* *i* **by**
blast

sublocale *real-sigma-finite-filtered-measure* ⊆ *sigma-finite-subalgebra* *M* *F* |*i*| **by**
fastforce

5.3 Finite Filtered Measure

locale *finite-filtered-measure* = *filtered-measure* + *finite-measure*

sublocale *finite-filtered-measure* ⊆ *sigma-finite-filtered-measure*
using *subalgebras* **by** (*unfold-locale*, *blast*, *meson dual-order.refl finite-measure-axioms finite-measure-def finite-measure-restr-to-subalg sigma-finite-measure.sigma-finite-countable*)

locale *nat-finite-filtered-measure* = *finite-filtered-measure* *M* *F* 0 :: *nat* **for** *M* *F*

locale *real-finite-filtered-measure* = *finite-filtered-measure* *M* *F* 0 :: *real* **for** *M* *F*

sublocale *nat-finite-filtered-measure* ⊆ *nat-sigma-finite-filtered-measure* ..

sublocale *real-finite-filtered-measure* ⊆ *real-sigma-finite-filtered-measure* ..

5.4 Constant Filtration

lemma *filtered-measure-constant-filtration*:

assumes *subalgebra* *M* *F*

shows *filtered-measure* *M* (λ-. *F*) *t*₀

using *assms* **by** (*unfold-locale*) *blast*+

sublocale *sigma-finite-subalgebra* ⊆ *constant-filtration*: *sigma-finite-filtered-measure*
M λ- :: 't :: {*second-countable-topology*, *linorder-topology*}. *F* *t*₀
using *subalg* **by** (*unfold-locale*) *blast*+

lemma (**in** *finite-measure*) *filtered-measure-constant-filtration*:

```

assumes subalgebra  $M$   $F$ 
shows finite-filtered-measure  $M$   $(\lambda \cdot. F)$   $t_0$ 
using assms by (unfold-locales) blast+

end

```

```

theory Stochastic-Process
imports Filtered-Measure Measure-Space-Supplement HOL-Probability.Independent-Family
begin

```

6 Stochastic Processes

6.1 Stochastic Process

A stochastic process is a collection of random variables, indexed by a type $'b$.

```

locale stochastic-process =
  fixes  $M$   $t_0$  and  $X :: 'b :: \{second-countable-topology, order-topology, t2-space\} \Rightarrow$ 
 $'a \Rightarrow 'c :: \{second-countable-topology, banach\}$ 
  assumes random-variable[measurable]:  $\bigwedge i. t_0 \leq i \implies X\ i \in borel-measurable\ M$ 
begin

```

```

definition left-continuous where left-continuous =  $(\lambda E\ \xi\ in\ M. \forall t. continuous$ 
 $(at-left\ t)\ (\lambda i. X\ i\ \xi))$ 

```

```

definition right-continuous where right-continuous =  $(\lambda E\ \xi\ in\ M. \forall t. continuous$ 
 $(at-right\ t)\ (\lambda i. X\ i\ \xi))$ 

```

```

end

```

```

locale nat-stochastic-process = stochastic-process  $M$   $0 :: nat$   $X$  for  $M$   $X$ 
locale real-stochastic-process = stochastic-process  $M$   $0 :: real$   $X$  for  $M$   $X$ 

```

```

lemma stochastic-process-const-fun:
  assumes  $f \in borel-measurable\ M$ 
  shows stochastic-process  $M$   $t_0$   $(\lambda \cdot. f)$  using assms by (unfold-locales)

```

```

lemma stochastic-process-const:
  shows stochastic-process  $M$   $t_0$   $(\lambda i \cdot. c\ i)$  by (unfold-locales) simp

```

```

context stochastic-process
begin

```

```

lemma compose-stochastic:
  assumes  $\bigwedge i. t_0 \leq i \implies f\ i \in borel-measurable\ borel$ 
  shows stochastic-process  $M$   $t_0$   $(\lambda i\ \xi. (f\ i)\ (X\ i\ \xi))$ 
  by (unfold-locales) (intro measurable-compose[OF random-variable assms])

```


lemma *norm-stochastic*: *stochastic-process* $M\ t_0\ (\lambda i\ \xi.\ \text{norm}\ (X\ i\ \xi))$ **by** (*fastforce*
intro: *compose-stochastic*)

lemma *scaleR-right-stochastic*:
assumes *stochastic-process* $M\ t_0\ Y$
shows *stochastic-process* $M\ t_0\ (\lambda i\ \xi.\ (Y\ i\ \xi) *_{\mathbb{R}} (X\ i\ \xi))$
using *stochastic-process.random-variable*[*OF assms*] *random-variable* **by** (*unfold-locales*)
simp

lemma *scaleR-right-const-fun-stochastic*:
assumes $f \in \text{borel-measurable}\ M$
shows *stochastic-process* $M\ t_0\ (\lambda i\ \xi.\ f\ \xi *_{\mathbb{R}} (X\ i\ \xi))$
by (*unfold-locales*) (*intro borel-measurable-scaleR assms random-variable*)

lemma *scaleR-right-const-stochastic*: *stochastic-process* $M\ t_0\ (\lambda i\ \xi.\ c\ i *_{\mathbb{R}} (X\ i\ \xi))$
by (*unfold-locales*) *simp*

lemma *add-stochastic*:
assumes *stochastic-process* $M\ t_0\ Y$
shows *stochastic-process* $M\ t_0\ (\lambda i\ \xi.\ X\ i\ \xi + Y\ i\ \xi)$
using *stochastic-process.random-variable*[*OF assms*] *random-variable* **by** (*unfold-locales*)
simp

lemma *diff-stochastic*:
assumes *stochastic-process* $M\ t_0\ Y$
shows *stochastic-process* $M\ t_0\ (\lambda i\ \xi.\ X\ i\ \xi - Y\ i\ \xi)$
using *stochastic-process.random-variable*[*OF assms*] *random-variable* **by** (*unfold-locales*)
simp

lemma *uminus-stochastic*: *stochastic-process* $M\ t_0\ (-X)$ **using** *scaleR-right-const-stochastic*[*of*
 $\lambda -. -1]$ **by** (*simp add: fun-Compl-def*)

lemma *partial-sum-stochastic*: *stochastic-process* $M\ t_0\ (\lambda n\ \xi.\ \sum_{i \in \{t_0..n\}} X\ i\ \xi)$
by (*unfold-locales*) *simp*

lemma *partial-sum'-stochastic*: *stochastic-process* $M\ t_0\ (\lambda n\ \xi.\ \sum_{i \in \{t_0..<n\}} X\ i\ \xi)$
by (*unfold-locales*) *simp*

end

lemma *stochastic-process-sum*:
assumes $\bigwedge i.\ i \in I \implies \text{stochastic-process}\ M\ t_0\ (X\ i)$
shows *stochastic-process* $M\ t_0\ (\lambda k\ \xi.\ \sum_{i \in I} X\ i\ k\ \xi)$ **using** *assms*[*THEN*
stochastic-process.random-variable] **by** (*unfold-locales, auto*)

6.1.1 Natural Filtration

The natural filtration induced by a stochastic process X is the filtration generated by all events involving the process up to the time index t , i.e. Σ

$t = \sigma \{X \ s \mid s. s \leq t\}.$

definition *natural-filtration* :: 'a measure \Rightarrow 'b \Rightarrow ('b \Rightarrow 'a \Rightarrow 'c :: topological-space) \Rightarrow 'b :: {second-countable-topology, order-topology} \Rightarrow 'a measure **where**
natural-filtration $M \ t_0 \ Y = (\lambda t. \text{family-vimage-algebra } (\text{space } M) \ \{Y \ i \mid i. i \in \{t_0..t\}\} \text{ borel})$

abbreviation *nat-natural-filtration* $\equiv \lambda M. \text{natural-filtration } M \ (0 :: \text{nat})$

abbreviation *real-natural-filtration* $\equiv \lambda M. \text{natural-filtration } M \ (0 :: \text{real})$

lemma *space-natural-filtration[simp]*: $\text{space } (\text{natural-filtration } M \ t_0 \ X \ t) = \text{space } M$ **unfolding** *natural-filtration-def* *space-family-vimage-algebra* ..

lemma *sets-natural-filtration*: $\text{sets } (\text{natural-filtration } M \ t_0 \ X \ t) = \text{sigma-sets } (\text{space } M) \ (\bigcup i \in \{t_0..t\}. \{X \ i - 'A \cap \text{space } M \mid A. A \in \text{borel}\})$

unfolding *natural-filtration-def* *sets-family-vimage-algebra* **by** (*intro sigma-sets-eqI*) *blast+*

lemma *sets-natural-filtration'*:

assumes *borel* = *sigma UNIV S*

shows $\text{sets } (\text{natural-filtration } M \ t_0 \ X \ t) = \text{sigma-sets } (\text{space } M) \ (\bigcup i \in \{t_0..t\}. \{X \ i - 'A \cap \text{space } M \mid A. A \in S\})$

proof (*subst sets-natural-filtration, intro sigma-sets-eqI, clarify*)

fix *i* **and** *A* :: 'a set **assume** *asm*: $i \in \{t_0..t\} \ A \in \text{sets borel}$

hence $A \in \text{sigma-sets UNIV S}$ **unfolding** *assms* **by** *simp*

thus $X \ i - 'A \cap \text{space } M \in \text{sigma-sets } (\text{space } M) \ (\bigcup i \in \{t_0..t\}. \{X \ i - 'A \cap \text{space } M \mid A. A \in S\})$

proof (*induction*)

case (*Compl a*)

have $X \ i - '(UNIV - a) \cap \text{space } M = \text{space } M - (X \ i - 'a \cap \text{space } M)$ **by** *blast*

then show ?*case* **using** *Compl(2)[THEN sigma-sets.Compl]* **by** *presburger*

next

case (*Union a*)

have $X \ i - '\bigcup (\text{range } a) \cap \text{space } M = \bigcup (\text{range } (\lambda j. X \ i - 'a \ j \cap \text{space } M))$

by *blast*

then show ?*case* **using** *Union(2)[THEN sigma-sets.Union]* **by** *presburger*

qed (*auto intro: asm sigma-sets.Empty*)

qed (*intro sigma-sets.Basic, force simp add: assms*)

lemma *sets-natural-filtration-open*:

$\text{sets } (\text{natural-filtration } M \ t_0 \ X \ t) = \text{sigma-sets } (\text{space } M) \ (\bigcup i \in \{t_0..t\}. \{X \ i - 'A \cap \text{space } M \mid A. \text{open } A\})$

using *sets-natural-filtration'* **by** (*force simp only: borel-def mem-Collect-eq*)

lemma *sets-natural-filtration-oi*:

$\text{sets } (\text{natural-filtration } M \ t_0 \ X \ t) = \text{sigma-sets } (\text{space } M) \ (\bigcup i \in \{t_0..t\}. \{X \ i - 'A \cap \text{space } M \mid A :: - :: \{\text{linorder-topology, second-countable-topology}\} \text{ set. } A \in \text{range greaterThan}\})$

by (*rule sets-natural-filtration'[OF borel-Ioi]*)

lemma *sets-natural-filtration-io:*

*sets (natural-filtration M t₀ X t) = sigma-sets (space M) (⋃ i ∈ {t₀..t}. {X i - ' A
 ∩ space M | A :: - :: {linorder-topology, second-countable-topology} set. A ∈ range
 lessThan})*

by (rule sets-natural-filtration'[OF borel-Iio])

lemma *sets-natural-filtration-ci:*

*sets (natural-filtration M t₀ X t) = sigma-sets (space M) (⋃ i ∈ {t₀..t}. {X i - '
 A ∩ space M | A :: real set. A ∈ range atLeast})*

by (rule sets-natural-filtration'[OF borel-Ici])

context *stochastic-process*

begin

lemma *subalgebra-natural-filtration:*

shows *subalgebra M (natural-filtration M t₀ X i)*

unfolding *subalgebra-def* **using** *measurable-family-iff-sets* **by** (*force simp add:*
natural-filtration-def)

lemma *filtered-measure-natural-filtration:*

shows *filtered-measure M (natural-filtration M t₀ X) t₀*

by (*unfold-locales*) (*intro subalgebra-natural-filtration, simp only: sets-natural-filtration,*
intro sigma-sets-subseteq, force)

In order to show that the natural filtration constitutes a filtered sigma finite
 measure, we need to provide a countable exhausting set in the preimage of
 X t₀.

lemma *sigma-finite-filtered-measure-natural-filtration:*

assumes *exhausting-set: countable A (⋃ A) = space M ∧ a. a ∈ A ⇒ emeasure*
M a ≠ ∞ ∧ a. a ∈ A ⇒ ∃ b ∈ borel. a = X t₀ - ' b ∩ space M

shows *sigma-finite-filtered-measure M (natural-filtration M t₀ X) t₀*

proof (*unfold-locales*)

have *A ⊆ sets (restr-to-subalg M (natural-filtration M t₀ X t₀))* **using** *exhaust-*
ing-set **by** (*simp add: sets-restr-to-subalg[OF subalgebra-natural-filtration] sets-natural-filtration*)
fast

moreover have *⋃ A = space (restr-to-subalg M (natural-filtration M t₀ X t₀))*

unfolding *space-restr-to-subalg* **using** *exhausting-set* **by** *simp*

moreover have *∀ a ∈ A. emeasure (restr-to-subalg M (natural-filtration M t₀ X*
t₀)) a ≠ ∞ **using** *calculation(1) exhausting-set(3)*

by (*auto simp add: sets-restr-to-subalg[OF subalgebra-natural-filtration] emea-*
sure-restr-to-subalg[OF subalgebra-natural-filtration])

ultimately show *∃ A. countable A ∧ A ⊆ sets (restr-to-subalg M (natural-filtration*
M t₀ X t₀)) ∧ ⋃ A = space (restr-to-subalg M (natural-filtration M t₀ X t₀)) ∧
(∀ a ∈ A. emeasure (restr-to-subalg M (natural-filtration M t₀ X t₀)) a ≠ ∞) **using**
exhausting-set **by** *blast*

show *⋀ i j. [t₀ ≤ i; i ≤ j] ⇒ sets (natural-filtration M t₀ X i) ⊆ sets (natural-filtration*
M t₀ X j) **using** *filtered-measure.subalgebra-F[OF filtered-measure-natural-filtration]*
by (*simp add: subalgebra-def*)

qed (*auto intro: subalgebra-natural-filtration*)

lemma *finite-filtered-measure-natural-filtration*:

assumes *finite-measure M*

shows *finite-filtered-measure M (natural-filtration M t₀ X) t₀*

using *finite-measure.axioms[OF assms] filtered-measure-natural-filtration by intro-locales*

end

Filtration generated by independent variables.

lemma (*in prob-space*) *indep-set-natural-filtration*:

assumes $t_0 \leq s < t$ *indep-vars ($\lambda \cdot$. borel) X {t₀..}*

shows *indep-set (natural-filtration M t₀ X s) (vimage-algebra (space M) (X t) borel)*

proof –

have *indep-sets (λi . {X i – ‘A \cap space M | A. A \in sets borel}) (\bigcup (range (case-bool {t₀..s} {t})))*

using *assms*

by (*intro assms(3)[unfolded indep-vars-def, THEN conjunct2, THEN indep-sets-mono]*)
(*auto simp add: case-bool-if*)

thus *?thesis unfolding indep-set-def using assms*

by (*intro indep-sets-cong[THEN iffD1, OF refl - indep-sets-collect-sigma[of λi . {X i – ‘A \cap space M | A. A \in borel} case-bool {t₀..s} {t}]]]*)

(*simp add: sets-natural-filtration sets-vimage-algebra split: bool.split, simp, intro Int-stableI, clarsimp,metis sets.Int vimage-Int Int-commute Int-left-absorb Int-left-commute, force simp add: disjoint-family-on-def split: bool.split*)

qed

6.2 Adapted Process

We call a collection a stochastic process X adapted if $X\ i$ is $F\ i$ -borel-measurable for all indices i .

locale *adapted-process* = *filtered-measure M F t₀* **for** $M\ F\ t_0$ **and** $X :: - \Rightarrow - \Rightarrow - :: \{second-countable-topology, banach\} +$

assumes *adapted[measurable]: $\bigwedge i. t_0 \leq i \implies X\ i \in borel-measurable\ (F\ i)$*

begin

lemma *adaptedE[elim]*:

assumes $\llbracket \bigwedge i. t_0 \leq j \implies j \leq i \implies X\ j \in borel-measurable\ (F\ i) \rrbracket \implies P$

shows P

using *assms using adapted by (metis dual-order.trans borel-measurable-subalgebra sets-F-mono space-F)*

lemma *adaptedD*:

assumes $t_0 \leq j < i$

shows $X\ j \in borel-measurable\ (F\ i)$ **using** *assms adaptedE by meson*

end

locale *nat-adapted-process* = *adapted-process* $M\ F\ 0 :: \text{nat } X$ **for** $M\ F\ X$
locale *real-adapted-process* = *adapted-process* $M\ F\ 0 :: \text{real } X$ **for** $M\ F\ X$

sublocale *nat-adapted-process* \subseteq *nat-filtered-measure* ..
sublocale *real-adapted-process* \subseteq *real-filtered-measure* ..

lemma (in *filtered-measure*) *adapted-process-const-fun*:
assumes $f \in \text{borel-measurable } (F\ t_0)$
shows *adapted-process* $M\ F\ t_0\ (\lambda\ -. \ f)$
using *measurable-from-subalg subalgebra-F assms* **by** (*unfold-locales*) *blast*

lemma (in *filtered-measure*) *adapted-process-const*:
shows *adapted-process* $M\ F\ t_0\ (\lambda i\ -. \ c\ i)$ **by** (*unfold-locales*) *simp*

context *adapted-process*
begin

lemma *compose-adapted*:
assumes $\bigwedge i. t_0 \leq i \implies f\ i \in \text{borel-measurable borel}$
shows *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. \ (f\ i)\ (X\ i\ \xi))$
by (*unfold-locales*) (*intro measurable-compose[OF adapted assms]*)

lemma *norm-adapted*: *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. \ \text{norm } (X\ i\ \xi))$ **by** (*fastforce intro: compose-adapted*)

lemma *scaleR-right-adapted*:
assumes *adapted-process* $M\ F\ t_0\ R$
shows *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. \ (R\ i\ \xi) *_{\mathbb{R}} (X\ i\ \xi))$
using *adapted-process.adapted[OF assms]* *adapted* **by** (*unfold-locales*) *simp*

lemma *scaleR-right-const-fun-adapted*:
assumes $f \in \text{borel-measurable } (F\ t_0)$
shows *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. \ f\ \xi *_{\mathbb{R}} (X\ i\ \xi))$
using *assms* **by** (*fast intro: scaleR-right-adapted adapted-process-const-fun*)

lemma *scaleR-right-const-adapted*: *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. \ c\ i *_{\mathbb{R}} (X\ i\ \xi))$
by (*unfold-locales*) *simp*

lemma *add-adapted*:
assumes *adapted-process* $M\ F\ t_0\ Y$
shows *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. \ X\ i\ \xi + Y\ i\ \xi)$
using *adapted-process.adapted[OF assms]* *adapted* **by** (*unfold-locales*) *simp*

lemma *diff-adapted*:
assumes *adapted-process* $M\ F\ t_0\ Y$
shows *adapted-process* $M\ F\ t_0\ (\lambda i\ \xi. \ X\ i\ \xi - Y\ i\ \xi)$
using *adapted-process.adapted[OF assms]* *adapted* **by** (*unfold-locales*) *simp*

lemma *uminus-adapted*: *adapted-process* $M F t_0 (-X)$ **using** *scaleR-right-const-adapted*[*of* $\lambda \cdot, -1$] **by** (*simp add: fun-Compl-def*)

lemma *partial-sum-adapted*: *adapted-process* $M F t_0 (\lambda n \xi. \sum_{i \in \{t_0..n\}} X i \xi)$
proof (*unfold-locales*)
 fix $i :: 'b$
 have $X j \in \text{borel-measurable } (F i)$ **if** $t_0 \leq j$ **for** j **using** *that adaptedE* **by** *meson*
 thus $(\lambda \xi. \sum_{i \in \{t_0..i\}} X i \xi) \in \text{borel-measurable } (F i)$ **by** *simp*
qed

lemma *partial-sum'-adapted*: *adapted-process* $M F t_0 (\lambda n \xi. \sum_{i \in \{t_0..<n\}} X i \xi)$

proof (*unfold-locales*)
 fix $i :: 'b$
 have $X j \in \text{borel-measurable } (F i)$ **if** $t_0 \leq j$ **for** $j < i$ **using** *that adaptedE* **by** *fastforce*
 thus $(\lambda \xi. \sum_{i \in \{t_0..<i\}} X i \xi) \in \text{borel-measurable } (F i)$ **by** *simp*
qed

end

lemma (*in nat-adapted-process*) *partial-sum-Suc-adapted*: *nat-adapted-process* $M F (\lambda n \xi. \sum_{i < n} X (Suc i) \xi)$
proof (*unfold-locales*)
 fix i
 have $X j \in \text{borel-measurable } (F i)$ **if** $j \leq i$ **for** j **using** *that adaptedD* **by** *blast*
 thus $(\lambda \xi. \sum_{i < i} X (Suc i) \xi) \in \text{borel-measurable } (F i)$ **by** *auto*
qed

lemma (*in filtered-measure*) *adapted-process-sum*:
 assumes $\bigwedge i. i \in I \implies \text{adapted-process } M F t_0 (X i)$
 shows *adapted-process* $M F t_0 (\lambda k \xi. \sum_{i \in I} X i k \xi)$
proof –
 {
 fix $i k$ **assume** $i \in I$ **and** *asm*: $t_0 \leq k$
then interpret *adapted-process* $M F t_0 X i$ **using** *assms* **by** *simp*
 have $X i k \in \text{borel-measurable } M X i k \in \text{borel-measurable } (F k)$ **using** *measurable-from-subalg subalgebras adapted asm* **by** (*blast, simp*)
 }
 thus *?thesis* **by** (*unfold-locales*) *simp*
qed

An adapted process is necessarily a stochastic process.

sublocale *adapted-process* \subseteq *stochastic-process* **using** *measurable-from-subalg subalgebras adapted* **by** (*unfold-locales*) *blast*

sublocale *nat-adapted-process* \subseteq *nat-stochastic-process* ..

sublocale *real-adapted-process* \subseteq *real-stochastic-process* ..

A stochastic process is always adapted to the natural filtration it generates.

lemma (*in stochastic-process*) *adapted-process-natural-filtration*: *adapted-process*
M (natural-filtration M t₀ X) t₀ X
using *filtered-measure-natural-filtration*
by (*intro-locales*) (*auto simp add: natural-filtration-def intro!:* *adapted-process-axioms.intro*
measurable-family-vimage-algebra)

6.3 Progressively Measurable Process

locale *progressive-process* = *filtered-measure M F t₀* **for** *M F t₀* **and** *X :: - \Rightarrow -*
 \Rightarrow *- :: {second-countable-topology, banach} +*
assumes *progressive[measurable]: $\bigwedge t. t_0 \leq t \Rightarrow (\lambda(i, x). X i x) \in \text{borel-measurable}$*
(restrict-space borel {t₀..t}) \otimes_M F t)
begin

lemma *progressiveD*:
assumes *S \in borel*
shows *($\lambda(j, \xi). X j \xi$) - ' S \cap ({t₀..i} \times space M) \in (restrict-space borel {t₀..i}*
 \otimes_M F i)
using *measurable-sets[OF progressive, OF - assms, of i]*
by (*cases t₀ \leq i*) (*auto simp add: space-restrict-space sets-pair-measure space-pair-measure*)

end

locale *nat-progressive-process* = *progressive-process M F 0 :: nat X* **for** *M F X*
locale *real-progressive-process* = *progressive-process M F 0 :: real X* **for** *M F X*

lemma (*in filtered-measure*) *progressive-process-const-fun*:
assumes *f \in borel-measurable (F t₀)*
shows *progressive-process M F t₀ (λ -. f)*
proof (*unfold-locales*)
fix *i* **assume** *asm: t₀ \leq i*
have *f \in borel-measurable (F i)* **using** *borel-measurable-mono[OF order.refl asm]*
assms **by** *blast*
thus *case-prod (λ -. f) \in borel-measurable (restrict-space borel {t₀..i} \otimes_M F i)*
using *measurable-compose[OF measurable-snd]* **by** *simp*
qed

lemma (*in filtered-measure*) *progressive-process-const*:
assumes *c \in borel-measurable borel*
shows *progressive-process M F t₀ (λ i -. c i)*
using *assms* **by** (*unfold-locales*) (*auto simp add: measurable-split-conv intro!:*
measurable-compose[OF measurable-fst] measurable-restrict-space1)

context *progressive-process*
begin

lemma *compose-progressive*:

assumes *case-prod* $f \in \text{borel-measurable borel}$

shows *progressive-process* $M F t_0 (\lambda i \xi. (f i) (X i \xi))$

proof

fix i **assume** *asm*: $t_0 \leq i$

have $(\lambda(j, \xi). (j, X j \xi)) \in (\text{restrict-space borel } \{t_0..i\} \otimes_M F i) \rightarrow_M \text{borel } \otimes_M \text{borel}$

using *progressive*[*OF asm*] *measurable-fst*''[*OF measurable-restrict-space1*, *OF measurable-id*]

by (*auto simp add: measurable-pair-iff measurable-split-conv*)

moreover have $(\lambda(j, \xi). f j (X j \xi)) = \text{case-prod } f \circ ((\lambda(j, y). (j, y)) \circ (\lambda(j, \xi). (j, X j \xi)))$ **by** *fastforce*

ultimately show $(\lambda(j, \xi). (f j) (X j \xi)) \in \text{borel-measurable } (\text{restrict-space borel } \{t_0..i\} \otimes_M F i)$ **using** *assms* **by** (*simp add: borel-prod*)

qed

lemma *norm-progressive*: *progressive-process* $M F t_0 (\lambda i \xi. \text{norm } (X i \xi))$ **using** *measurable-compose*[*OF progressive borel-measurable-norm*] **by** (*unfold-locales*) *simp*

lemma *scaleR-right-progressive*:

assumes *progressive-process* $M F t_0 R$

shows *progressive-process* $M F t_0 (\lambda i \xi. (R i \xi) *_{\mathbb{R}} (X i \xi))$

using *progressive-process.progressive*[*OF assms*] **by** (*unfold-locales*) (*simp add: progressive assms*)

lemma *scaleR-right-const-fun-progressive*:

assumes $f \in \text{borel-measurable } (F t_0)$

shows *progressive-process* $M F t_0 (\lambda i \xi. f \xi *_{\mathbb{R}} (X i \xi))$

using *assms* **by** (*fast intro: scaleR-right-progressive progressive-process-const-fun*)

lemma *scaleR-right-const-progressive*:

assumes $c \in \text{borel-measurable borel}$

shows *progressive-process* $M F t_0 (\lambda i \xi. c i *_{\mathbb{R}} (X i \xi))$

using *assms* **by** (*fastforce intro: scaleR-right-progressive progressive-process-const*)

lemma *add-progressive*:

assumes *progressive-process* $M F t_0 Y$

shows *progressive-process* $M F t_0 (\lambda i \xi. X i \xi + Y i \xi)$

using *progressive-process.progressive*[*OF assms*] **by** (*unfold-locales*) (*simp add: progressive assms*)

lemma *diff-progressive*:

assumes *progressive-process* $M F t_0 Y$

shows *progressive-process* $M F t_0 (\lambda i \xi. X i \xi - Y i \xi)$

using *progressive-process.progressive*[*OF assms*] **by** (*unfold-locales*) (*simp add: progressive assms*)

lemma *uminus-progressive*: *progressive-process* $M F t_0 (-X)$ **using** *scaleR-right-const-progressive*[*of*

$\lambda -. -1]$ **by** (*simp add: fun-Compl-def*)

end

A progressively measurable process is also adapted.

sublocale *progressive-process* \subseteq *adapted-process* **using** *measurable-compose-rev*[*OF progressive measurable-Pair1* \uparrow]
unfolding *prod.case space-restrict-space*
by *unfold-locales simp*

sublocale *nat-progressive-process* \subseteq *nat-adapted-process* ..

sublocale *real-progressive-process* \subseteq *real-adapted-process* ..

In the discrete setting, adaptedness is equivalent to progressive measurability.

theorem *nat-progressive-iff-adapted*: *nat-progressive-process* $M\ F\ X \longleftrightarrow$ *nat-adapted-process* $M\ F\ X$

proof (*intro iffI*)

assume *asm*: *nat-progressive-process* $M\ F\ X$

interpret *nat-progressive-process* $M\ F\ X$ **by** (*rule asm*)

show *nat-adapted-process* $M\ F\ X$..

next

assume *asm*: *nat-adapted-process* $M\ F\ X$

interpret *nat-adapted-process* $M\ F\ X$ **by** (*rule asm*)

show *nat-progressive-process* $M\ F\ X$

proof (*unfold-locales, intro borel-measurableI*)

fix $S :: 'b\ set$ **and** $i :: nat$ **assume** *open-S*: *open S*

{

fix j **assume** *asm*: $j \leq i$

hence $X\ j - 'S \cap space\ M \in F\ i$ **using** *adaptedD*[*of j, THEN measurable-sets*]
space-F open-S **by** *fastforce*

moreover have *case-prod* $X - 'S \cap \{j\} \times space\ M = \{j\} \times (X\ j - 'S \cap space\ M)$ **for** j **by** *fast*

moreover have $\{j :: nat\} \in restrict\ space\ borel\ \{0..i\}$ **using** *asm* **by** (*simp add: sets-restrict-space-iff*)

ultimately have *case-prod* $X - 'S \cap \{j\} \times space\ M \in restrict\ space\ borel\ \{0..i\} \otimes_M F\ i$ **by** *simp*

}

hence $(\lambda j. (\lambda(x, y). X\ x\ y) - 'S \cap \{j\} \times space\ M) \cdot \{0..i\} \subseteq restrict\ space\ borel\ \{0..i\} \otimes_M F\ i$ **by** *blast*

moreover have *case-prod* $X - 'S \cap space\ (restrict\ space\ borel\ \{0..i\} \otimes_M F\ i) = (\bigcup_{j \leq i}. case\ prod\ X - 'S \cap \{j\} \times space\ M)$ **unfolding** *space-pair-measure space-restrict-space space-F* **by** *force*

ultimately show *case-prod* $X - 'S \cap space\ (restrict\ space\ borel\ \{0..i\} \otimes_M F\ i) \in restrict\ space\ borel\ \{0..i\} \otimes_M F\ i$ **by** (*metis sets.countable-UN*)

qed

qed

6.4 Predictable Process

We introduce the constant Σ_P to denote the predictable sigma algebra.

context *linearly-filtered-measure*
begin

definition $\Sigma_P :: ('b \times 'a) \text{ measure where predictable-sigma: } \Sigma_P \equiv \text{sigma } (\{t_0..\} \times \text{space } M) (\{\{s<..t\} \times A \mid A \text{ s t. } A \in F \text{ s } \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in F \text{ } t_0\})$

lemma *space-predictable-sigma[simp]: space $\Sigma_P = (\{t_0..\} \times \text{space } M)$ unfolding predictable-sigma space-measure-of-conv by blast*

lemma *sets-predictable-sigma: sets $\Sigma_P = \text{sigma-sets } (\{t_0..\} \times \text{space } M) (\{\{s<..t\} \times A \mid A \text{ s t. } A \in F \text{ s } \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in F \text{ } t_0\})$ unfolding predictable-sigma using space-F sets.sets-into-space by (subst sets-measure-of) fastforce+*

lemma *measurable-predictable-sigma-snd:*

assumes *countable \mathcal{I} $\mathcal{I} \subseteq \{\{s<..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\} \{t_0<..\} \subseteq (\bigcup \mathcal{I})$*

shows *$\text{snd} \in \Sigma_P \rightarrow_M F \text{ } t_0$*

proof *(intro measurableI)*

fix *$S :: 'a \text{ set assume } \text{asm: } S \in F \text{ } t_0$*

have *countable: countable $((\lambda I. I \times S) \text{ ' } \mathcal{I})$ using $\text{assms}(1)$ by blast*

have *$(\lambda I. I \times S) \text{ ' } \mathcal{I} \subseteq \{\{s<..t\} \times A \mid A \text{ s t. } A \in F \text{ s } \wedge t_0 \leq s \wedge s < t\}$ using $\text{sets-F-mono}[OF \text{ order-refl, THEN subsetD, OF - asm}] \text{ assms}(2)$ by blast*

hence *$(\bigcup I \in \mathcal{I}. I \times S) \cup \{t_0\} \times S \in \Sigma_P$ unfolding sets-predictable-sigma using asm by (intro sigma-sets-Un[OF sigma-sets-UNION[OF countable] sigma-sets.Basic] sigma-sets.Basic) blast+*

moreover have *$\text{snd} - 'S \cap \text{space } \Sigma_P = \{t_0..\} \times S$ using $\text{sets.sets-into-space}[OF \text{ asm}]$ by fastforce*

moreover have *$\{t_0\} \cup \{t_0<..\} = \{t_0..\}$ by auto*

moreover have *$(\bigcup I \in \mathcal{I}. I \times S) \cup \{t_0\} \times S = \{t_0..\} \times S$ using $\text{assms}(2,3)$ calculation(3) by fastforce*

ultimately show *$\text{snd} - 'S \cap \text{space } \Sigma_P \in \Sigma_P$ by argo*

qed *(auto)*

lemma *measurable-predictable-sigma-fst:*

assumes *countable \mathcal{I} $\mathcal{I} \subseteq \{\{s<..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\} \{t_0<..\} \subseteq (\bigcup \mathcal{I})$*

shows *$\text{fst} \in \Sigma_P \rightarrow_M \text{borel}$*

proof *–*

have *$A \times \text{space } M \in \text{sets } \Sigma_P$ if $A \in \text{sigma-sets } \{t_0..\} \{\{s<..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\}$ for A unfolding sets-predictable-sigma using that*

proof *(induction rule: sigma-sets.induct)*

case *(Basic a)*

thus *?case using space-F sets.top by blast*

next

case *(Compl a)*

have *$(\{t_0..\} - a) \times \text{space } M = \{t_0..\} \times \text{space } M - a \times \text{space } M$ by blast*

```

    then show ?case using Compl(2)[THEN sigma-sets.Compl] by presburger
next
  case (Union a)
  have  $\bigcup (range\ a) \times space\ M = \bigcup (range\ (\lambda i. a\ i \times space\ M))$  by blast
  then show ?case using Union(2)[THEN sigma-sets.Union] by presburger
qed (auto)
moreover have restrict-space borel  $\{t_0..\} = sigma\ \{t_0..\}\ \{\{s<..t\} \mid s\ t. t_0 \leq s$ 
 $\wedge s < t\}$ 
proof -
  have sigma-sets  $\{t_0..\}\ ((\cap)\ \{t_0..\}\ 'sigma-sets\ UNIV\ (range\ greaterThan)) =$ 
sigma-sets  $\{t_0..\}\ \{\{s<..t\} \mid s\ t. t_0 \leq s \wedge s < t\}$ 
proof (intro sigma-sets-eqI ; clarify)
  fix A :: 'b set assume asm:  $A \in sigma-sets\ UNIV\ (range\ greaterThan)$ 
  thus  $\{t_0..\} \cap A \in sigma-sets\ \{t_0..\}\ \{\{s<..t\} \mid s\ t. t_0 \leq s \wedge s < t\}$ 
  proof (induction rule: sigma-sets.induct)
    case (Basic a)
    then obtain s where  $a = \{s<..\}$  by blast
    show ?case
    proof (cases  $t_0 \leq s$ )
      case True
      hence *:  $\{t_0..\} \cap a = (\bigcup i \in \mathcal{I}. \{s<..\} \cap i)$  using s assms(3) by force
      have  $((\cap)\ \{s<..\}\ ' \mathcal{I}) \subseteq sigma-sets\ \{t_0..\}\ \{\{s<..t\} \mid s\ t. t_0 \leq s \wedge s < t\}$ 
      proof (clarify)
        fix A assume  $A \in \mathcal{I}$ 
        then obtain s' t' where  $A: A = \{s'<..t'\}$   $t_0 \leq s' s' < t'$  using assms(2)
      by blast
      hence  $\{s<..\} \cap A = \{max\ s\ s'<..t'\}$  by fastforce
      moreover have  $t_0 \leq max\ s\ s'$  using A True by linarith
      moreover have  $max\ s\ s' < t'$  if  $s < t'$  using A that by linarith
      moreover have  $\{s<..\} \cap A = \{\}$  if  $\neg s < t'$  using A that by force
      ultimately show  $\{s<..\} \cap A \in sigma-sets\ \{t_0..\}\ \{\{s<..t\} \mid s\ t. t_0 \leq s$ 
 $\wedge s < t\}$  by (cases  $s < t'$ ) (blast, simp add: sigma-sets.Empty)
    qed
    thus ?thesis unfolding * using assms(1) by (intro sigma-sets-UNION)
  auto
next
  case False
  hence  $\{t_0..\} \cap a = \{t_0..\}$  using s by force
  thus ?thesis using sigma-sets-top by auto
qed
next
  case (Compl a)
  have  $\{t_0..\} \cap (UNIV - a) = \{t_0..\} - (\{t_0..\} \cap a)$  by blast
  then show ?case using Compl(2)[THEN sigma-sets.Compl] by presburger
next
  case (Union a)
  have  $\{t_0..\} \cap \bigcup (range\ a) = \bigcup (range\ (\lambda i. \{t_0..\} \cap a\ i))$  by blast
  then show ?case using Union(2)[THEN sigma-sets.Union] by presburger
qed (simp add: sigma-sets.Empty)

```

```

next
  fix s t assume asm:  $t_0 \leq s < t$ 
  hence *:  $\{s <..t\} = \{s <..\} \cap (\{t_0..\} - \{t <..\})$  by force
  have  $\{s <..\} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\})$  'sigma-sets UNIV (range greaterThan))
using asm by (intro sigma-sets.Basic) auto
  moreover have  $\{t_0..\} - \{t <..\} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\})$  'sigma-sets
UNIV (range greaterThan)) using asm by (intro sigma-sets.Compl sigma-sets.Basic)
auto
  ultimately show  $\{s <..t\} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\})$  'sigma-sets UNIV
(range greaterThan)) unfolding * Int-range-binary[of  $\{s <..\}$ ] by (intro sigma-sets-Inter[OF
- binary-in-sigma-sets]) auto
qed
  thus ?thesis unfolding borel-Ioi restrict-space-def emeasure-sigma by (force
intro: sigma-eqI)
qed
ultimately have restrict-space borel  $\{t_0..\} \otimes_M \text{sigma } (\text{space } M) \{\} \subseteq \text{sets } \Sigma_P$ 

  unfolding sets-pair-measure space-restrict-space space-measure-of-conv
  using space-predictable-sigma sets.sigma-algebra-axioms[of  $\Sigma_P$ ]
  by (intro sigma-algebra.sigma-sets-subset) (auto simp add: sigma-sets-empty-eq
sets-measure-of-conv)
  moreover have space (restrict-space borel  $\{t_0..\} \otimes_M \text{sigma } (\text{space } M) \{\}) =$ 
space  $\Sigma_P$  by (simp add: space-pair-measure)
  moreover have  $\text{fst} \in \text{restrict-space borel } \{t_0..\} \otimes_M \text{sigma } (\text{space } M) \{\} \rightarrow_M$ 
borel by (fastforce intro: measurable-fst''[OF measurable-restrict-space1, of  $\lambda x. x$ ])

  ultimately show ?thesis by (meson borel-measurable-subalgebra)
qed
end

locale predictable-process = linearly-filtered-measure M F  $t_0$  for M F  $t_0$  and X ::
-  $\Rightarrow$  -  $\Rightarrow$  - :: {second-countable-topology, banach} +
  assumes predictable:  $(\lambda(t, x). X t x) \in \text{borel-measurable } \Sigma_P$ 
begin

lemmas predictableD = measurable-sets[OF predictable, unfolded space-predictable-sigma]

end

locale nat-predictable-process = predictable-process M F 0 :: nat X for M F X
locale real-predictable-process = predictable-process M F 0 :: real X for M F X

lemma (in nat-filtered-measure) measurable-predictable-sigma-snd':
  shows  $\text{snd} \in \Sigma_P \rightarrow_M F 0$ 
  by (intro measurable-predictable-sigma-snd[of range  $(\lambda x. \{\text{Suc } x\})$ ]) (force | simp
add: greaterThan-0)+

```

lemma (in *nat-filtered-measure*) *measurable-predictable-sigma-fst'*:
shows $\text{fst} \in \Sigma_P \rightarrow_M \text{borel}$
by (intro *measurable-predictable-sigma-fst*[of range $(\lambda x. \{\text{Suc } x\})$]) (force | simp
add: greaterThan-0)+

lemma (in *real-filtered-measure*) *measurable-predictable-sigma-snd'*:
shows $\text{snd} \in \Sigma_P \rightarrow_M F\ 0$
using *real-arch-simple* **by** (intro *measurable-predictable-sigma-snd*[of range $(\lambda x::\text{nat}.$
 $\{0 < ..\text{real } (\text{Suc } x)\})$]) (fastforce intro: *add-increasing*)+

lemma (in *real-filtered-measure*) *measurable-predictable-sigma-fst'*:
shows $\text{fst} \in \Sigma_P \rightarrow_M \text{borel}$
using *real-arch-simple* **by** (intro *measurable-predictable-sigma-fst*[of range $(\lambda x::\text{nat}.$
 $\{0 < ..\text{real } (\text{Suc } x)\})$]) (fastforce intro: *add-increasing*)+

lemma (in *linearly-filtered-measure*) *predictable-process-const-fun*:
assumes $\text{snd} \in \Sigma_P \rightarrow_M F\ t_0$ $f \in \text{borel-measurable } (F\ t_0)$
shows *predictable-process* $M\ F\ t_0\ (\lambda\cdot. f)$
using *measurable-compose-rev*[*OF assms*(2)] *assms*(1) **by** (*unfold-locale*) (*auto*
simp add: measurable-split-conv)

lemma (in *nat-filtered-measure*) *predictable-process-const-fun'*[*intro*]:
assumes $f \in \text{borel-measurable } (F\ 0)$
shows *nat-predictable-process* $M\ F\ (\lambda\cdot. f)$
using *assms* **by** (intro *predictable-process-const-fun*[*OF measurable-predictable-sigma-snd'*,
THEN nat-predictable-process.intro])

lemma (in *real-filtered-measure*) *predictable-process-const-fun'*[*intro*]:
assumes $f \in \text{borel-measurable } (F\ 0)$
shows *real-predictable-process* $M\ F\ (\lambda\cdot. f)$
using *assms* **by** (intro *predictable-process-const-fun*[*OF measurable-predictable-sigma-snd'*,
THEN real-predictable-process.intro])

lemma (in *linearly-filtered-measure*) *predictable-process-const*:
assumes $\text{fst} \in \text{borel-measurable } \Sigma_P$ $c \in \text{borel-measurable borel}$
shows *predictable-process* $M\ F\ t_0\ (\lambda i\cdot. c\ i)$
using *assms* **by** (*unfold-locale*) (*simp add: measurable-split-conv*)

lemma (in *linearly-filtered-measure*) *predictable-process-const-const*[*intro*]:
shows *predictable-process* $M\ F\ t_0\ (\lambda\cdot\cdot. c)$
by (*unfold-locale*) *simp*

lemma (in *nat-filtered-measure*) *predictable-process-const'*[*intro*]:
assumes $c \in \text{borel-measurable borel}$
shows *nat-predictable-process* $M\ F\ (\lambda i\cdot. c\ i)$
using *assms* **by** (intro *predictable-process-const*[*OF measurable-predictable-sigma-fst'*,

THEN nat-predictable-process.intro])

lemma (in *real-filtered-measure*) *predictable-process-const'*[*intro*]:
 assumes $c \in \text{borel-measurable borel}$
 shows *real-predictable-process* $M F (\lambda i \cdot c i)$
 using *assms* by (intro *predictable-process-const*[*OF measurable-predictable-sigma-fst'*],
THEN real-predictable-process.intro)

context *predictable-process*
begin

lemma *compose-predictable*:
 assumes $\text{fst} \in \text{borel-measurable } \Sigma_P \text{ case-prod } f \in \text{borel-measurable borel}$
 shows *predictable-process* $M F t_0 (\lambda i \xi. (f i) (X i \xi))$
proof
 have $(\lambda(i, \xi). (i, X i \xi)) \in \Sigma_P \rightarrow_M \text{borel} \otimes_M \text{borel}$ using *predictable assms(1)*
 by (auto simp add: *measurable-pair-iff measurable-split-conv*)
 moreover have $(\lambda(i, \xi). f i (X i \xi)) = \text{case-prod } f o (\lambda(i, \xi). (i, X i \xi))$ by
fastforce
 ultimately show $(\lambda(i, \xi). f i (X i \xi)) \in \text{borel-measurable } \Sigma_P$ unfolding *borel-prod*
 using *assms* by *simp*
qed

lemma *norm-predictable*: *predictable-process* $M F t_0 (\lambda i \xi. \text{norm } (X i \xi))$ using
measurable-compose[*OF predictable borel-measurable-norm*]
 by (unfold-locale) (simp add: *prod.case-distrib*)

lemma *scaleR-right-predictable*:
 assumes *predictable-process* $M F t_0 R$
 shows *predictable-process* $M F t_0 (\lambda i \xi. (R i \xi) *_R (X i \xi))$
 using *predictable predictable-process.predictable*[*OF assms*] by (unfold-locale)
 (auto simp add: *measurable-split-conv*)

lemma *scaleR-right-const-fun-predictable*:
 assumes $\text{snd} \in \Sigma_P \rightarrow_M F t_0 f \in \text{borel-measurable } (F t_0)$
 shows *predictable-process* $M F t_0 (\lambda i \xi. f \xi *_R (X i \xi))$
 using *assms* by (fast intro: *scaleR-right-predictable predictable-process-const-fun*)

lemma *scaleR-right-const-predictable*:
 assumes $\text{fst} \in \text{borel-measurable } \Sigma_P c \in \text{borel-measurable borel}$
 shows *predictable-process* $M F t_0 (\lambda i \xi. c i *_R (X i \xi))$
 using *assms* by (fastforce intro: *scaleR-right-predictable predictable-process-const*)

lemma *scaleR-right-const'-predictable*: *predictable-process* $M F t_0 (\lambda i \xi. c *_R (X i \xi))$
 by (fastforce intro: *scaleR-right-predictable*)

lemma *add-predictable*:
 assumes *predictable-process* $M F t_0 Y$

shows *predictable-process* $M F t_0 (\lambda i \xi. X i \xi + Y i \xi)$
using *predictable* *predictable-process.predictable*[*OF assms*] **by** (*unfold-locales*)
(auto simp add: measurable-split-conv)

lemma *diff-predictable*:

assumes *predictable-process* $M F t_0 Y$
shows *predictable-process* $M F t_0 (\lambda i \xi. X i \xi - Y i \xi)$
using *predictable* *predictable-process.predictable*[*OF assms*] **by** (*unfold-locales*)
(auto simp add: measurable-split-conv)

lemma *uminus-predictable*: *predictable-process* $M F t_0 (-X)$ **using** *scaleR-right-const'-predictable*[*of -1*] **by** (*simp add: fun-Comp-def*)

end

Every predictable process is also progressively measurable.

sublocale *predictable-process* \subseteq *progressive-process*

proof (*unfold-locales*)

fix $i :: 'b$ **assume** $asm: t_0 \leq i$
 $\{$
fix $S :: ('b \times 'a)$ *set* **assume** $S \in \{\{s <..t\} \times A \mid A s t. A \in F s \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in F t_0\}$
hence $(\lambda x. x) - ' S \cap (\{t_0..i\} \times \text{space } M) \in \text{restrict-space borel } \{t_0..i\} \otimes_M F$
 i

proof

assume $S \in \{\{s <..t\} \times A \mid A s t. A \in F s \wedge t_0 \leq s \wedge s < t\}$
then obtain $s t A$ **where** $S\text{-is}: S = \{s <..t\} \times A$ $t_0 \leq s$ $s < t$ $A \in F s$ **by**
blast

hence $(\lambda x. x) - ' S \cap (\{t_0..i\} \times \text{space } M) = \{s <.. \min i t\} \times A$ **using**
sets.sets-into-space[*OF S-is(4)*] **by** *auto*

then show *?thesis* **using** $S\text{-is}$ *sets-F-mono*[*of s i*] **by** (*cases s ≤ i*) (*fastforce simp add: sets-restrict-space-iff*) +

next

assume $S \in \{\{t_0\} \times A \mid A. A \in F t_0\}$
then obtain A **where** $S\text{-is}: S = \{t_0\} \times A$ $A \in F t_0$ **by** *blast*
hence $(\lambda x. x) - ' S \cap (\{t_0..i\} \times \text{space } M) = \{t_0\} \times A$ **using** asm *sets.sets-into-space*[*OF S-is(2)*] **by** *auto*

thus *?thesis* **using** $S\text{-is}(2)$ *sets-F-mono*[*OF order-refl asm*] asm **by** (*fastforce simp add: sets-restrict-space-iff*)

qed

hence $(\lambda x. x) - ' S \cap \text{space } (\text{restrict-space borel } \{t_0..i\} \otimes_M F i) \in \text{restrict-space borel } \{t_0..i\} \otimes_M F i$ **by** (*simp add: space-pair-measure space-F*[*OF asm*])

$\}$

moreover have $\{\{s <..t\} \times A \mid A s t. A \in \text{sets } (F s) \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\} \times A \mid A. A \in \text{sets } (F t_0)\} \subseteq \text{Pow } (\{t_0..i\} \times \text{space } M)$ **using** *sets.sets-into-space* **by**
force

ultimately have $(\lambda x. x) \in \text{restrict-space borel } \{t_0..i\} \otimes_M F i \rightarrow_M \Sigma_P$ **using** *space-F*[*OF asm*] **by** (*intro measurable-sigma-sets*[*OF sets-predictable-sigma*])
(fast, force simp add: space-pair-measure)

thus $\text{case-prod } X \in \text{borel-measurable } (\text{restrict-space borel } \{t_0..i\} \otimes_M F i)$ **using**
predictable by simp
qed

sublocale $\text{nat-predictable-process} \subseteq \text{nat-progressive-process} ..$
sublocale $\text{real-predictable-process} \subseteq \text{real-progressive-process} ..$

The following lemma characterizes predictability in a discrete-time setting.

lemma (in *nat-filtered-measure*) *sets-in-filtration*:

assumes $(\bigcup i. \{i\} \times A i) \in \Sigma_P$
shows $A (Suc i) \in F i \ A \ 0 \in F \ 0$
using *assms unfolding sets-predictable-sigma*
proof (*induction* $(\bigcup i. \{i\} \times A i)$ *arbitrary: A*)
case *Basic*
 {
assume $\exists S. (\bigcup i. \{i\} \times A i) = \{0\} \times S$
then obtain S **where** $S: (\bigcup i. \{i\} \times A i) = \{bot\} \times S$ **unfolding** *bot-nat-def*
by *blast*
hence $S \in F \ bot$ **using** *Basic* **by** (*fastforce simp add: times-eq-iff bot-nat-def*)
moreover have $A \ i = \{\}$ **if** $i \neq bot$ **for** i **using** *that S* **by** *blast*
moreover have $A \ bot = S$ **using** S **by** *blast*
ultimately have $A (Suc i) \in F i \ A \ 0 \in F \ 0$ **for** i **unfolding** *bot-nat-def* **by**
(auto simp add: bot-nat-def)
 }
note $*$ **=** *this*
 {
assume $\nexists S. (\bigcup i. \{i\} \times A i) = \{0\} \times S$
then obtain $s \ t \ B$ **where** $B: (\bigcup i. \{i\} \times A i) = \{s<..t\} \times B \ B \in \text{sets } (F \ s)$
 $s < t$ **using** *Basic* **by** *auto*
hence $A \ i = B$ **if** $i \in \{s<..t\}$ **for** i **using** *that* **by** *fast*
moreover have $A \ i = \{\}$ **if** $i \notin \{s<..t\}$ **for** i **using** B **that** **by** *fastforce*
ultimately have $A (Suc i) \in F i \ A \ 0 \in F \ 0$ **for** i **unfolding** *bot-nat-def* **using**
 B *sets-F-mono* **by** (*auto simp add: bot-nat-def*) (*metis less-Suc-eq-le sets.empty-sets subset-eq*)
 }
note $**$ **=** *this*
show $A (Suc i) \in \text{sets } (F i) \ A \ 0 \in \text{sets } (F \ 0)$ **using** $*(1)[of \ i] \ *(2) \ ***(1)[of \ i]$
 $**(2)$ **by** *blast+*
next
case *Empty*
 {
case *1*
then show *?case* **using** *Empty* **by** *simp*
next
case *2*
then show *?case* **using** *Empty* **by** *simp*
 }
next
case (*Compl a*)


```

have  $a\text{-in}: a \subseteq \{0..\} \times \text{space } M$  using  $\text{Compl}(1)$  sets.sets-into-space sets-predictable-sigma
space-predictable-sigma by metis
  hence  $A\text{-in}: A \ i \subseteq \text{space } M$  for  $i$  using  $\text{Compl}(4)$  by blast
  have  $a: a = \{0..\} \times \text{space } M - (\bigcup i. \{i\} \times A \ i)$  using  $a\text{-in}$   $\text{Compl}(4)$  by blast
  also have  $\dots = -(\bigcap j. -(\{j\} \times (\text{space } M - A \ j)))$  by blast
  also have  $\dots = (\bigcup j. \{j\} \times (\text{space } M - A \ j))$  by blast
  finally have  $*$ :  $(\text{space } M - A \ (\text{Suc } i)) \in F \ i \ (\text{space } M - A \ 0) \in F \ 0$  using
 $\text{Compl}(2,3)$  by auto
  {
    case 1
    then show  $?case$  using  $*$   $A\text{-in}$  by (metis bot-nat-0.extremum double-diff
sets.Diff sets.top sets-F-mono sets-le-imp-space-le space-F)
  next
    case 2
    then show  $?case$  using  $*$   $A\text{-in}$  by (metis bot-nat-0.extremum double-diff
sets.Diff sets.top sets-F-mono sets-le-imp-space-le space-F)
  }
next
  case ( $\text{Union } a$ )
  have  $a\text{-in}: a \ i \subseteq \{0..\} \times \text{space } M$  for  $i$  using  $\text{Union}(1)$  sets.sets-into-space
sets-predictable-sigma space-predictable-sigma by metis
  hence  $A\text{-in}: A \ i \subseteq \text{space } M$  for  $i$  using  $\text{Union}(4)$  by blast
  have  $\text{snd } x \in \text{snd } ' (a \ i \cap (\{fst \ x\} \times \text{space } M))$  if  $x \in a \ i$  for  $i \ x$  using that
a-in by fastforce
  hence  $a\text{-i}: a \ i = (\bigcup j. \{j\} \times (\text{snd } ' (a \ i \cap (\{j\} \times \text{space } M))))$  for  $i$  by force
  have  $A\text{-i}: A \ i = \text{snd } ' (\bigcup (\text{range } a) \cap (\{i\} \times \text{space } M))$  for  $i$  unfolding
 $\text{Union}(4)$  using  $A\text{-in}$  by force
  have  $*$ :  $\text{snd } ' (a \ j \cap (\{\text{Suc } i\} \times \text{space } M)) \in F \ i \ \text{snd } ' (a \ j \cap (\{0\} \times \text{space } M))$ 
 $\in F \ 0$  for  $j$  using  $\text{Union}(2,3)[OF \ a\text{-i}]$  by auto
  {
    case 1
    have  $(\bigcup j. \text{snd } ' (a \ j \cap (\{\text{Suc } i\} \times \text{space } M))) \in F \ i$  using  $*$  by fast
    moreover have  $(\bigcup j. \text{snd } ' (a \ j \cap (\{\text{Suc } i\} \times \text{space } M))) = \text{snd } ' (\bigcup (\text{range}$ 
a) \cap (\{\text{Suc } i\} \times \text{space } M)) by fast
    ultimately show  $?case$  using  $A\text{-i}$  by metis
  next
    case 2
    have  $(\bigcup j. \text{snd } ' (a \ j \cap (\{0\} \times \text{space } M))) \in F \ 0$  using  $*$  by fast
    moreover have  $(\bigcup j. \text{snd } ' (a \ j \cap (\{0\} \times \text{space } M))) = \text{snd } ' (\bigcup (\text{range } a) \cap$ 
 $(\{0\} \times \text{space } M))$  by fast
    ultimately show  $?case$  using  $A\text{-i}$  by metis
  }
}
qed

```

This leads to the following useful fact.

lemma (*in nat-predictable-process*) *adapted-Suc*: *nat-adapted-process* $M \ F \ (\lambda i. \ X$
 $(\text{Suc } i))$
proof (*unfold-locales, intro borel-measurableI*)
fix $S :: 'b \text{ set}$ **and** i **assume** *open-S*: *open* S

have $\{Suc\ i\} = \{i < .. Suc\ i\}$ **by** *fastforce*
hence $\{Suc\ i\} \times space\ M \in \Sigma_P$ **using** *space-F[symmetric, of i]* **unfolding**
sets-predictable-sigma **by** (*intro sigma-sets.Basic*) *blast*
moreover **have** $case\text{-}prod\ X - 'S \cap (UNIV \times space\ M) \in \Sigma_P$ **unfolding**
atLeast-0[symmetric] **using** *open-S* **by** (*intro predictableD, simp add: borel-open*)
ultimately **have** $case\text{-}prod\ X - 'S \cap (\{Suc\ i\} \times space\ M) \in \Sigma_P$ **unfolding**
sets-predictable-sigma **using** *space-F sets.sets-into-space*
by (*subst Times-Int-distrib1[of {Suc i} UNIV space M, simplified], subst*
inf commute, subst Int-assoc[symmetric], subst Int-range-binary)
(intro sigma-sets-Inter binary-in-sigma-sets, fast)+
moreover **have** $case\text{-}prod\ X - 'S \cap (\{Suc\ i\} \times space\ M) = \{Suc\ i\} \times (X\ (Suc\ i) - 'S \cap space\ M)$ **by** (*auto simp add: le-Suc-eq*)
moreover **have** $... = (\bigcup j. \{j\} \times (if\ j = Suc\ i\ then\ (X\ (Suc\ i) - 'S \cap space\ M)\ else\ \{\}))$ **by** (*force split: if-splits*)
ultimately **have** $(\bigcup j. \{j\} \times (if\ j = Suc\ i\ then\ (X\ (Suc\ i) - 'S \cap space\ M)\ else\ \{\})) \in \Sigma_P$ **by** *argo*
thus $X\ (Suc\ i) - 'S \cap space\ (F\ i) \in sets\ (F\ i)$ **using** *sets-in-filtration[of $\lambda j. if\ j = Suc\ i\ then\ (X\ (Suc\ i) - 'S \cap space\ M)\ else\ \{\}$ space-F[OF zero-le]* **by**
presburger
qed

theorem *nat-predictable-process-iff*: $nat\text{-}predictable\text{-}process\ M\ F\ X \longleftrightarrow nat\text{-}adapted\text{-}process\ M\ F\ (\lambda i. X\ (Suc\ i)) \wedge X\ 0 \in borel\text{-}measurable\ (F\ 0)$

proof (*intro iffI*)

assume *asm*: $nat\text{-}adapted\text{-}process\ M\ F\ (\lambda i. X\ (Suc\ i)) \wedge X\ 0 \in borel\text{-}measurable\ (F\ 0)$

interpret *nat-adapted-process* $M\ F\ \lambda i. X\ (Suc\ i)$ **using** *asm* **by** *blast*

have $(\lambda(x, y). X\ x\ y) \in borel\text{-}measurable\ \Sigma_P$

proof (*intro borel-measurableI*)

fix $S :: 'b\ set$ **assume** *open-S*: *open S*

have $\{i\} \times (X\ i - 'S \cap space\ M) \in sets\ \Sigma_P$ **for** i

proof (*cases i*)

case 0

then show *?thesis* **unfolding** *sets-predictable-sigma*

using *measurable-sets[OF - borel-open[OF open-S], of X 0 F 0]* *asm* **by** *auto*

next

case $(Suc\ i)$

have $\{Suc\ i\} = \{i < .. Suc\ i\}$ **by** *fastforce*

then show *?thesis* **unfolding** *sets-predictable-sigma*

using *measurable-sets[OF adapted borel-open[OF open-S], of i]*

by (*intro sigma-sets.Basic, auto simp add: Suc*)

qed

moreover **have** $(\lambda(x, y). X\ x\ y) - 'S \cap space\ \Sigma_P = (\bigcup i. \{i\} \times (X\ i - 'S \cap space\ M))$ **by** *fastforce*

ultimately show $(\lambda(x, y). X\ x\ y) - 'S \cap space\ \Sigma_P \in sets\ \Sigma_P$ **by** *simp*

qed

thus $nat\text{-}predictable\text{-}process\ M\ F\ X$ **by** (*unfold-locales*)

next

assume *asm*: $nat\text{-}predictable\text{-}process\ M\ F\ X$

```

interpret nat-predictable-process  $M\ F\ X$  by (rule asm)
show nat-adapted-process  $M\ F\ (\lambda i. X\ (Suc\ i)) \wedge X\ 0 \in \text{borel-measurable}\ (F\ 0)$ 
using adapted-Suc by simp
qed

end

```

```

theory Martingale
imports Stochastic-Process Conditional-Expectation-Banach
begin

```

7 Martingales

The following locales are necessary for defining martingales.

7.1 Additional Locale Definitions

```

locale sigma-finite-adapted-process = sigma-finite-filtered-measure  $M\ F\ t_0$  + adapted-process
 $M\ F\ t_0\ X$  for  $M\ F\ t_0\ X$ 

```

```

locale nat-sigma-finite-adapted-process = sigma-finite-adapted-process  $M\ F\ 0$  :: nat
 $X$  for  $M\ F\ X$ 

```

```

locale real-sigma-finite-adapted-process = sigma-finite-adapted-process  $M\ F\ 0$  ::
real  $X$  for  $M\ F\ X$ 

```

```

sublocale nat-sigma-finite-adapted-process  $\subseteq$  nat-sigma-finite-filtered-measure ..
sublocale real-sigma-finite-adapted-process  $\subseteq$  real-sigma-finite-filtered-measure ..

```

```

locale finite-adapted-process = finite-filtered-measure  $M\ F\ t_0$  + adapted-process  $M\ F\ t_0\ X$ 
for  $M\ F\ t_0\ X$ 

```

```

sublocale finite-adapted-process  $\subseteq$  sigma-finite-adapted-process ..

```

```

locale nat-finite-adapted-process = finite-adapted-process  $M\ F\ 0$  :: nat  $X$  for  $M\ F\ X$ 

```

```

locale real-finite-adapted-process = finite-adapted-process  $M\ F\ 0$  :: real  $X$  for  $M\ F\ X$ 

```

```

sublocale nat-finite-adapted-process  $\subseteq$  nat-sigma-finite-adapted-process ..
sublocale real-finite-adapted-process  $\subseteq$  real-sigma-finite-adapted-process ..

```

```

locale sigma-finite-adapted-process-order = sigma-finite-adapted-process  $M\ F\ t_0\ X$ 
for  $M\ F\ t_0$  and  $X$  :: -  $\Rightarrow$  -  $\Rightarrow$  - :: {order-topology, ordered-real-vector}

```

locale *nat-sigma-finite-adapted-process-order* = *sigma-finite-adapted-process-order*
M F 0 :: nat X for M F X
locale *real-sigma-finite-adapted-process-order* = *sigma-finite-adapted-process-order*
M F 0 :: real X for M F X

sublocale *nat-sigma-finite-adapted-process-order* \subseteq *nat-sigma-finite-adapted-process*
..
sublocale *real-sigma-finite-adapted-process-order* \subseteq *real-sigma-finite-adapted-process*
..

locale *finite-adapted-process-order* = *finite-adapted-process* *M F t₀ X for M F t₀*
and *X :: - \Rightarrow - \Rightarrow - :: {order-topology, ordered-real-vector}*

locale *nat-finite-adapted-process-order* = *finite-adapted-process-order* *M F 0 :: nat*
X for M F X
locale *real-finite-adapted-process-order* = *finite-adapted-process-order* *M F 0 :: real*
X for M F X

sublocale *nat-finite-adapted-process-order* \subseteq *nat-sigma-finite-adapted-process-order*
..
sublocale *real-finite-adapted-process-order* \subseteq *real-sigma-finite-adapted-process-order*
..

locale *sigma-finite-adapted-process-linorder* = *sigma-finite-adapted-process-order*
M F t₀ X for M F t₀ and X :: - \Rightarrow - \Rightarrow - :: {linorder-topology}

locale *nat-sigma-finite-adapted-process-linorder* = *sigma-finite-adapted-process-linorder*
M F 0 :: nat X for M F X
locale *real-sigma-finite-adapted-process-linorder* = *sigma-finite-adapted-process-linorder*
M F 0 :: real X for M F X

sublocale *nat-sigma-finite-adapted-process-linorder* \subseteq *nat-sigma-finite-adapted-process-order*
..
sublocale *real-sigma-finite-adapted-process-linorder* \subseteq *real-sigma-finite-adapted-process-order*
..

locale *finite-adapted-process-linorder* = *finite-adapted-process-order* *M F t₀ X for*
M F t₀ and X :: - \Rightarrow - \Rightarrow - :: {linorder-topology}

locale *nat-finite-adapted-process-linorder* = *finite-adapted-process-linorder* *M F 0*
:: nat X for M F X
locale *real-finite-adapted-process-linorder* = *finite-adapted-process-linorder* *M F 0*
:: real X for M F X

sublocale *nat-finite-adapted-process-linorder* \subseteq *nat-sigma-finite-adapted-process-linorder*
..
sublocale *real-finite-adapted-process-linorder* \subseteq *real-sigma-finite-adapted-process-linorder*

..

7.2 Martingale

locale *martingale* = *sigma-finite-adapted-process* +
assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
and *martingale-property*: $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE \xi \text{ in } M. X i \xi =$
cond-exp $M (F i) (X j) \xi$

locale *martingale-order* = *martingale* $M F t_0 X$ **for** $M F t_0$ **and** $X :: - \Rightarrow - \Rightarrow -$
 $:: \{ \text{order-topology, ordered-real-vector} \}$
locale *martingale-linorder* = *martingale* $M F t_0 X$ **for** $M F t_0$ **and** $X :: - \Rightarrow - \Rightarrow -$
 $- :: \{ \text{linorder-topology, ordered-real-vector} \}$
sublocale *martingale-linorder* \subseteq *martingale-order* ..

lemma (**in** *sigma-finite-filtered-measure*) *martingale-const-fun*[*intro*]:
assumes *integrable* $M f f \in \text{borel-measurable } (F t_0)$
shows *martingale* $M F t_0 (\lambda-. f)$
using *assms sigma-finite-subalgebra.cond-exp-F-meas*[*OF - assms*(1), *THEN AE-symmetric*]
borel-measurable-mono
by (*unfold-locale*) *blast*+

lemma (**in** *sigma-finite-filtered-measure*) *martingale-cond-exp*[*intro*]:
assumes *integrable* $M f$
shows *martingale* $M F t_0 (\lambda i. \text{cond-exp } M (F i) f)$
using *sigma-finite-subalgebra.borel-measurable-cond-exp'* *borel-measurable-cond-exp*

by (*unfold-locale*) (*auto intro: sigma-finite-subalgebra.cond-exp-nested-subalg*[*OF - assms*]
simp add: subalgebra-F subalgebras)

corollary (**in** *sigma-finite-filtered-measure*) *martingale-zero*[*intro*]: *martingale* M
 $F t_0 (\lambda-. 0)$ **by** *fastforce*

corollary (**in** *finite-filtered-measure*) *martingale-const*[*intro*]: *martingale* $M F t_0$
 $(\lambda-. c)$ **by** *fastforce*

7.3 Submartingale

locale *submartingale* = *sigma-finite-adapted-process-order* +
assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
and *submartingale-property*: $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE \xi \text{ in } M. X i \xi \leq$
cond-exp $M (F i) (X j) \xi$

locale *submartingale-linorder* = *submartingale* $M F t_0 X$ **for** $M F t_0$ **and** $X :: -$
 $\Rightarrow - \Rightarrow - :: \{ \text{linorder-topology} \}$

sublocale *martingale-order* \subseteq *submartingale* **using** *martingale-property* **by** (*unfold-locale*)
(*force simp add: integrable*)
sublocale *martingale-linorder* \subseteq *submartingale-linorder* ..

7.4 Supermartingale

locale *supermartingale* = *sigma-finite-adapted-process-order* +
assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M \ (X \ i)$
and *supermartingale-property*: $\bigwedge i \ j. t_0 \leq i \implies i \leq j \implies AE \ \xi \text{ in } M. \ X \ i \ \xi$
 $\geq \text{cond-exp } M \ (F \ i) \ (X \ j) \ \xi$

locale *supermartingale-linorder* = *supermartingale* *M F t₀ X* **for** *M F t₀* **and** *X*
 $:: - \Rightarrow - \Rightarrow - :: \{\text{linorder-topology}\}$

sublocale *martingale-order* \subseteq *supermartingale* **using** *martingale-property* **by** (*unfold-locales*)
(*force simp add: integrable*) +

sublocale *martingale-linorder* \subseteq *supermartingale-linorder* ..

A stochastic process is a martingale, if and only if it is both a submartingale and a supermartingale.

lemma *martingale-iff*:

shows *martingale* *M F t₀ X* \longleftrightarrow *submartingale* *M F t₀ X* \wedge *supermartingale* *M F t₀ X*

proof (*rule iffI*)

assume *asm*: *martingale* *M F t₀ X*

interpret *martingale-order* *M F t₀ X* **by** (*intro martingale-order.intro asm*)

show *submartingale* *M F t₀ X* \wedge *supermartingale* *M F t₀ X* **using** *submartingale-axioms* *supermartingale-axioms* **by** *blast*

next

assume *asm*: *submartingale* *M F t₀ X* \wedge *supermartingale* *M F t₀ X*

interpret *submartingale* *M F t₀ X* **by** (*simp add: asm*)

interpret *supermartingale* *M F t₀ X* **by** (*simp add: asm*)

show *martingale* *M F t₀ X* **using** *submartingale-property* *supermartingale-property*
by (*unfold-locales*) (*intro integrable, blast, force*)

qed

7.5 Martingale Lemmas

context *martingale*

begin

lemma *cond-exp-diff-eq-zero*:

assumes $t_0 \leq i \leq j$

shows $AE \ \xi \text{ in } M. \text{cond-exp } M \ (F \ i) \ (\lambda \xi. X \ j \ \xi - X \ i \ \xi) \ \xi = 0$

using *martingale-property*[*OF assms*] *assms*

sigma-finite-subalgebra.cond-exp-F-meas[*OF - integrable adapted, of i*]

sigma-finite-subalgebra.cond-exp-diff[*OF - integrable(1,1), of F i j i*] **by**

fastforce

lemma *set-integral-eq*:

assumes $A \in F \ i \ t_0 \leq i \leq j$

shows *set-lebesgue-integral* *M A* (*X i*) = *set-lebesgue-integral* *M A* (*X j*)

proof –

interpret *sigma-finite-subalgebra* *M F i* **using** *assms*(2) **by** *blast*

have $\int x \in A. X \ i \ x \ \partial M = \int x \in A. \text{cond-exp } M \ (F \ i) \ (X \ j) \ x \ \partial M$ **using**
martingale-property[*OF assms*(2,3)] *borel-measurable-cond-exp'* *assms subalgebras*
subalgebra-def **by** (*intro set-lebesgue-integral-cong-AE*[*OF - random-variable*]) *fast-*
force+
also have $\dots = \int x \in A. X \ j \ x \ \partial M$ **using** *assms* **by** (*auto simp: integrable intro:*
cond-exp-set-integral[symmetric])
finally show *?thesis* .
qed

lemma *scaleR-const*[*intro*]:
shows *martingale* $M \ F \ t_0 \ (\lambda i \ x. \ c *_{\mathbb{R}} X \ i \ x)$
proof –
{
 fix $i \ j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
 interpret *sigma-finite-subalgebra* $M \ F \ i$ **using** *asm* **by** *blast*
 have $AE \ x \ in \ M. \ c *_{\mathbb{R}} X \ i \ x = \text{cond-exp } M \ (F \ i) \ (\lambda x. \ c *_{\mathbb{R}} X \ j \ x)$ **us-**
ing *asm cond-exp-scaleR-right*[*OF integrable, of j, THEN AE-symmetric*] *martin-*
gale-property[*OF asm*] **by** *force*
}
thus *?thesis* **by** (*unfold-locales*) (*auto simp add: integrable martingale.integrable*)
qed

lemma *uminus*[*intro*]:
shows *martingale* $M \ F \ t_0 \ (- \ X)$
using *scaleR-const*[*of -1*] **by** (*force intro: back-subst*[*of martingale M F t_0*])

lemma *add*[*intro*]:
assumes *martingale* $M \ F \ t_0 \ Y$
shows *martingale* $M \ F \ t_0 \ (\lambda i \ \xi. \ X \ i \ \xi + Y \ i \ \xi)$
proof –
interpret $Y: \text{martingale } M \ F \ t_0 \ Y$ **by** (*rule assms*)
{
 fix $i \ j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
 hence $AE \ \xi \ in \ M. \ X \ i \ \xi + Y \ i \ \xi = \text{cond-exp } M \ (F \ i) \ (\lambda x. \ X \ j \ x + Y \ j \ x) \ \xi$
 using *sigma-finite-subalgebra.cond-exp-add*[*OF - integrable martingale.integrable*][*OF*
assms], *of F i j j, THEN AE-symmetric*] *martingale-property*[*OF asm*] *martingale.martingale-property*[*OF assms*
asm] **by** *force*
}
thus *?thesis* **using** *assms*
by (*unfold-locales*) (*auto simp add: integrable martingale.integrable*)
qed

lemma *diff*[*intro*]:
assumes *martingale* $M \ F \ t_0 \ Y$
shows *martingale* $M \ F \ t_0 \ (\lambda i \ x. \ X \ i \ x - Y \ i \ x)$
proof –
interpret $Y: \text{martingale } M \ F \ t_0 \ Y$ **by** (*rule assms*)
{

```

    fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
    hence AE  $\xi$  in  $M$ .  $X i \xi - Y i \xi = \text{cond-exp } M (F i) (\lambda x. X j x - Y j x) \xi$ 
    using sigma-finite-subalgebra.cond-exp-diff[OF - integrable martingale.integrable[OF
assms], of  $F i j j$ , THEN AE-symmetric]
    martingale-property[OF asm] martingale.martingale-property[OF assms
asm] by fastforce
  }
  thus ?thesis using assms by (unfold-locales) (auto simp add: integrable martin-
gale.integrable)
qed

```

end

```

lemma (in sigma-finite-adapted-process) martingale-of-cond-exp-diff-eq-zero:
  assumes integrable:  $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$ 
  and diff-zero:  $\bigwedge i j. t_0 \leq i \implies i \leq j \implies \text{AE } x \text{ in } M. \text{cond-exp } M (F i) (\lambda \xi. X j \xi - X i \xi) x = 0$ 
  shows martingale  $M F t_0 X$ 
proof
  {
    fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
    thus AE  $\xi$  in  $M$ .  $X i \xi = \text{cond-exp } M (F i) (X j) \xi$ 
    using diff-zero[OF asm] sigma-finite-subalgebra.cond-exp-diff[OF - integrable(1,1),
of  $F i j i$ ]
    sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of  $i$ ] by
fastforce
  }
qed (intro integrable)

```

```

lemma (in sigma-finite-adapted-process) martingale-of-set-integral-eq:
  assumes integrable:  $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$ 
  and  $\bigwedge A i j. t_0 \leq i \implies i \leq j \implies A \in F i \implies \text{set-lebesgue-integral } M A (X i) = \text{set-lebesgue-integral } M A (X j)$ 
  shows martingale  $M F t_0 X$ 
proof (unfold-locales)
  fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
  interpret sigma-finite-subalgebra  $M F i$  using asm by blast
  interpret r: sigma-finite-measure restr-to-subalg  $M (F i)$  by (simp add: sigma-fin-subalg)
  {
    fix A assume  $A \in \text{restr-to-subalg } M (F i)$ 
    hence *:  $A \in F i$  using sets-restr-to-subalg subalgebras asm by blast
    have set-lebesgue-integral (restr-to-subalg  $M (F i)) A (X i) = \text{set-lebesgue-integral } M A (X i)$  using * subalg asm by (auto simp: set-lebesgue-integral-def intro: integral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator)
    also have ... = set-lebesgue-integral  $M A (\text{cond-exp } M (F i) (X j))$  using *
assms(2)[OF asm] cond-exp-set-integral[OF integrable] asm by auto
    finally have set-lebesgue-integral (restr-to-subalg  $M (F i)) A (X i) = \text{set-lebesgue-integral } M A (X j)$ 
(restr-to-subalg  $M (F i)) A (\text{cond-exp } M (F i) (X j))$  using * subalg by (auto simp:
set-lebesgue-integral-def intro!: integral-subalgebra2[symmetric] borel-measurable-scaleR

```


borel-measurable-cond-exp borel-measurable-indicator)
 }
 hence $AE \xi$ in $restr\text{-}to\text{-}subalg\ M\ (F\ i)$. $X\ i\ \xi = cond\text{-}exp\ M\ (F\ i)\ (X\ j)\ \xi$ **using** asm **by** (intro $r.density\text{-}unique\text{-}banach$, auto intro: $integrable\text{-}in\text{-}subalg\ subalg\ borel\text{-}measurable\text{-}cond\text{-}exp\ integrable$)
 thus $AE \xi$ in M . $X\ i\ \xi = cond\text{-}exp\ M\ (F\ i)\ (X\ j)\ \xi$ **using** $AE\text{-}restr\text{-}to\text{-}subalg\ [OF\ subalg]$ **by** blast
 qed (simp add: $integrable$)

7.6 Submartingale Lemmas

context $submartingale$
 begin

lemma $cond\text{-}exp\text{-}diff\text{-}nonneg$:
 assumes $t_0 \leq i\ i \leq j$
 shows $AE\ x$ in M . $cond\text{-}exp\ M\ (F\ i)\ (\lambda\xi. X\ j\ \xi - X\ i\ \xi)\ x \geq 0$
using $submartingale\text{-}property\ [OF\ assms]$ $assms\ sigma\text{-}finite\text{-}subalgebra.cond\text{-}exp\text{-}diff\ [OF\ -\ integrable\ (1,1),\ of\ -\ j\ i]\ sigma\text{-}finite\text{-}subalgebra.cond\text{-}exp\text{-}F\text{-}meas\ [OF\ -\ integrable\ adapted,\ of\ i]$ **by** fastforce

lemma $add\ [intro]$:
 assumes $submartingale\ M\ F\ t_0\ Y$
 shows $submartingale\ M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi + Y\ i\ \xi)$
proof –
 interpret Y : $submartingale\ M\ F\ t_0\ Y$ **by** (rule $assms$)
 {
 fix $i\ j :: 'b$ **assume** $asm: t_0 \leq i\ i \leq j$
 hence $AE\ \xi$ in M . $X\ i\ \xi + Y\ i\ \xi \leq cond\text{-}exp\ M\ (F\ i)\ (\lambda x. X\ j\ x + Y\ j\ x)\ \xi$
using $sigma\text{-}finite\text{-}subalgebra.cond\text{-}exp\text{-}add\ [OF\ -\ integrable\ submartingale.integrable\ [OF\ assms],\ of\ F\ i\ j\ j]$
 $submartingale\text{-}property\ [OF\ asm]\ submartingale.submartingale\text{-}property\ [OF\ assms\ asm]\ add\text{-}mono\ [of\ X\ i\ -\ -\ Y\ i\ -]$ **by** force
 }
 thus ?thesis **using** $assms$ **by** (unfold-locale) (auto simp add: $borel\text{-}measurable\text{-}add\ random\text{-}variable\ adapted\ integrable\ Y.random\text{-}variable\ Y.adapted\ submartingale.integrable$)

qed

lemma $diff\ [intro]$:
 assumes $supermartingale\ M\ F\ t_0\ Y$
 shows $supermartingale\ M\ F\ t_0\ (\lambda i\ \xi. X\ i\ \xi - Y\ i\ \xi)$
proof –
 interpret Y : $supermartingale\ M\ F\ t_0\ Y$ **by** (rule $assms$)
 {
 fix $i\ j :: 'b$ **assume** $asm: t_0 \leq i\ i \leq j$
 hence $AE\ \xi$ in M . $X\ i\ \xi - Y\ i\ \xi \leq cond\text{-}exp\ M\ (F\ i)\ (\lambda x. X\ j\ x - Y\ j\ x)\ \xi$
using $sigma\text{-}finite\text{-}subalgebra.cond\text{-}exp\text{-}diff\ [OF\ -\ integrable\ supermartingale.integrable\ [OF\ assms],\ of\ F\ i\ j\ j]$

```

      submartingale-property[OF asm] supermartingale.supermartingale-property[OF
assms asm] diff-mono[of  $X$   $i$  - - -  $Y$   $i$  -] by force
    }
    thus ?thesis using assms by (unfold-locale) (auto simp add: borel-measurable-diff
random-variable adapted integrable  $Y$ .random-variable  $Y$ .adapted supermartingale.integrable)

qed

```

```

lemma scaleR-nonneg:
  assumes  $c \geq 0$ 
  shows submartingale  $M$   $F$   $t_0$  ( $\lambda i \xi. c *_{\mathbb{R}} X i \xi$ )
proof
  {
    fix  $i j :: 'b$  assume asm:  $t_0 \leq i \leq j$ 
    thus  $\mathbb{A}E \xi$  in  $M. c *_{\mathbb{R}} X i \xi \leq \text{cond-exp } M (F i) (\lambda \xi. c *_{\mathbb{R}} X j \xi) \xi$ 
      using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of  $F i$ 
 $j$   $c$ ] submartingale-property[OF asm] by (fastforce intro!: scaleR-left-mono[OF -
assms])
  }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)

```

```

lemma scaleR-le-zero:
  assumes  $c \leq 0$ 
  shows supermartingale  $M$   $F$   $t_0$  ( $\lambda i \xi. c *_{\mathbb{R}} X i \xi$ )
proof
  {
    fix  $i j :: 'b$  assume asm:  $t_0 \leq i \leq j$ 
    thus  $\mathbb{A}E \xi$  in  $M. c *_{\mathbb{R}} X i \xi \geq \text{cond-exp } M (F i) (\lambda \xi. c *_{\mathbb{R}} X j \xi) \xi$ 
      using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of  $F i$   $j$ 
 $c$ ] submartingale-property[OF asm]
      by (fastforce intro!: scaleR-left-mono-neg[OF - assms])
  }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)

```

```

lemma uminus[intro]:
  shows supermartingale  $M$   $F$   $t_0$  ( $- X$ )
  unfolding fun-Comp-def using scaleR-le-zero[of  $-1$ ] by simp

end

```

```

context submartingale-linorder
begin

```

```

lemma set-integral-le:
  assumes  $A \in F$   $i t_0 \leq i \leq j$ 
  shows set-lebesgue-integral  $M$   $A$  ( $X i$ )  $\leq$  set-lebesgue-integral  $M$   $A$  ( $X j$ )
  using submartingale-property[OF assms(2), of  $j$ ] assms subalgebras

```

by (*subst sigma-finite-subalgebra.cond-exp-set-integral*[*OF* - *integrable assms*(1), of *j*])
 (*auto intro!*: *scaleR-left-mono integral-mono-AE-banach integrable-mult-indicator integrable simp add: subalgebra-def set-lebesgue-integral-def*)

lemma *max*:

assumes *submartingale-linorder* *M F t₀ Y*
shows *submartingale-linorder* *M F t₀ (λ i ξ. max (X i ξ) (Y i ξ))*
proof (*unfold-locales*)
interpret *Y*: *submartingale-linorder* *M F t₀ Y* **by** (*rule assms*)
 {
fix *i j* :: 'b **assume** *asm*: *t₀ ≤ i i ≤ j*
have *AE ξ in M. max (X i ξ) (Y i ξ) ≤ max (cond-exp M (F i) (X j) ξ)*
(cond-exp M (F i) (Y j) ξ) **using** *submartingale-property Y.submartingale-property*
asm unfolding max-def by fastforce
thus *AE ξ in M. max (X i ξ) (Y i ξ) ≤ cond-exp M (F i) (λ ξ. max (X j ξ) (Y*
j ξ)) ξ **using** *sigma-finite-subalgebra.cond-exp-max*[*OF* - *integrable Y.integrable, of*
F i j j] *asm by (fast intro: order.trans)*
 }
show $\bigwedge i. t_0 \leq i \implies (\lambda \xi. \max (X i \xi) (Y i \xi)) \in \text{borel-measurable } (F i) \bigwedge i.$
 $t_0 \leq i \implies \text{integrable } M (\lambda \xi. \max (X i \xi) (Y i \xi))$ **by** (*force intro: Y.integrable*
integrable assms)
qed

lemma *max-0*:

shows *submartingale-linorder* *M F t₀ (λ i ξ. max 0 (X i ξ))*
proof –
interpret *zero*: *martingale-linorder* *M F t₀ λ- -. 0* **by** (*force intro: martin-*
gale-linorder.intro martingale-order.intro)
show *?thesis* **by** (*intro zero.max submartingale-linorder.intro submartingale-axioms*)
qed

end

lemma (*in sigma-finite-adapted-process-order*) *submartingale-of-cond-exp-diff-nonneg*:

assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$
and *diff-nonneg*: $\bigwedge i j. t_0 \leq i \implies i \leq j \implies \text{AE } x \text{ in } M. \text{cond-exp } M (F i)$
 $(\lambda \xi. X j \xi - X i \xi) x \geq 0$
shows *submartingale* *M F t₀ X*
proof (*unfold-locales*)
 {
fix *i j* :: 'b **assume** *asm*: *t₀ ≤ i i ≤ j*
thus *AE ξ in M. X i ξ ≤ cond-exp M (F i) (X j) ξ*
using *diff-nonneg*[*OF asm*] *sigma-finite-subalgebra.cond-exp-diff*[*OF* - *inte-*
grable(1,1), of F i j i]
sigma-finite-subalgebra.cond-exp-F-meas[*OF* - *integrable adapted, of i*] **by**
fastforce
 }
qed (*intro integrable*)

```

lemma (in sigma-finite-adapted-process-linorder) submartingale-of-set-integral-le:
  assumes integrable:  $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$ 
    and  $\bigwedge A i j. t_0 \leq i \implies i \leq j \implies A \in F i \implies \text{set-lebesgue-integral } M A (X i) \leq \text{set-lebesgue-integral } M A (X j)$ 
  shows submartingale  $M F t_0 X$ 
proof (unfold-locales)
  {
    fix  $i j :: 'b$  assume  $asm: t_0 \leq i \leq j$ 
    interpret  $r: \text{sigma-finite-measure restr-to-subalg } M (F i)$  using asm sigma-finite-subalgebra.sigma-fin-subalg
  by blast
  {
    fix  $A$  assume  $A \in \text{restr-to-subalg } M (F i)$ 
    hence  $*$ :  $A \in F i$  using asm sets-restr-to-subalg subalgebras by blast
    have  $\text{set-lebesgue-integral } (\text{restr-to-subalg } M (F i)) A (X i) = \text{set-lebesgue-integral } M A (X i)$  using  $*$  asm subalgebras by (auto simp: set-lebesgue-integral-def intro: integral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator)
    also have  $\dots \leq \text{set-lebesgue-integral } M A (\text{cond-exp } M (F i) (X j))$  using  $*$  assms(2)[OF asm] asm sigma-finite-subalgebra.cond-exp-set-integral[OF - integrable] by fastforce
    also have  $\dots = \text{set-lebesgue-integral } (\text{restr-to-subalg } M (F i)) A (\text{cond-exp } M (F i) (X j))$  using  $*$  asm subalgebras by (auto simp: set-lebesgue-integral-def intro!: integral-subalgebra2[symmetric] borel-measurable-scaleR borel-measurable-cond-exp borel-measurable-indicator)
    finally have  $0 \leq \text{set-lebesgue-integral } (\text{restr-to-subalg } M (F i)) A (\lambda \xi. \text{cond-exp } M (F i) (X j) \xi - X i \xi)$  using  $*$  asm subalgebras by (subst set-integral-diff, auto simp add: set-integrable-def sets-restr-to-subalg intro!: integrable adapted integrable-in-subalg borel-measurable-scaleR borel-measurable-indicator borel-measurable-cond-exp integrable-mult-indicator)
  }
  hence  $AE \xi \text{ in } \text{restr-to-subalg } M (F i). 0 \leq \text{cond-exp } M (F i) (X j) \xi - X i \xi$ 
  by (intro r.density-nonneg integrable-in-subalg asm subalgebras borel-measurable-diff borel-measurable-cond-exp adapted Bochner-Integration.integrable-diff integrable-cond-exp integrable)
  thus  $AE \xi \text{ in } M. X i \xi \leq \text{cond-exp } M (F i) (X j) \xi$  using AE-restr-to-subalg[OF subalgebras] asm by simp
  }
qed (intro integrable)

```

7.7 Supermartingale Lemmas

The following lemmas are exact duals of the ones for submartingales.

```

context supermartingale
begin

```

```

lemma cond-exp-diff-nonneg:
  assumes  $t_0 \leq i \leq j$ 
  shows  $AE x \text{ in } M. \text{cond-exp } M (F i) (\lambda \xi. X i \xi - X j \xi) x \geq 0$ 
  using assms supermartingale-property[OF assms] sigma-finite-subalgebra.cond-exp-diff[OF

```

- *integrable*(1,1), of *F i i j*]
 sigma-finite-subalgebra.cond-exp-F-meas[*OF - integrable adapted*, of *i*] **by**
fastforce

lemma *add*[*intro*]:
 assumes *supermartingale M F t₀ Y*
 shows *supermartingale M F t₀ (λ i ξ. X i ξ + Y i ξ)*
proof -
 interpret *Y: supermartingale M F t₀ Y* **by** (*rule assms*)
 {
 fix *i j :: 'b* **assume** *asm: t₀ ≤ i i ≤ j*
 hence *AE ξ in M. X i ξ + Y i ξ ≥ cond-exp M (F i) (λ x. X j x + Y j x) ξ*
 using *sigma-finite-subalgebra.cond-exp-add*[*OF - integrable supermartingale.integrable*[*OF*
assms], of *F i j j*]
 supermartingale-property[*OF asm*] *supermartingale.supermartingale-property*[*OF*
assms asm] *add-mono*[of - *X i - - Y i -*] **by** *force*
 }
 thus *?thesis* **using** *assms* **by** (*unfold-locale*) (*auto simp add: borel-measurable-add*
random-variable adapted integrable Y.random-variable Y.adapted supermartingale.integrable)

qed

lemma *diff*[*intro*]:
 assumes *submartingale M F t₀ Y*
 shows *supermartingale M F t₀ (λ i ξ. X i ξ - Y i ξ)*
proof -
 interpret *Y: submartingale M F t₀ Y* **by** (*rule assms*)
 {
 fix *i j :: 'b* **assume** *asm: t₀ ≤ i i ≤ j*
 hence *AE ξ in M. X i ξ - Y i ξ ≥ cond-exp M (F i) (λ x. X j x - Y j x) ξ*
 using *sigma-finite-subalgebra.cond-exp-diff*[*OF - integrable submartingale.integrable*[*OF*
assms], of *F i j j*, *unfolded fun-diff-def*]
 supermartingale-property[*OF asm*] *submartingale.submartingale-property*[*OF*
assms asm] *diff-mono*[of - *X i - Y i -*] **by** *force*
 }
 thus *?thesis* **using** *assms* **by** (*unfold-locale*) (*auto simp add: borel-measurable-diff*
random-variable adapted integrable Y.random-variable Y.adapted submartingale.integrable)

qed

lemma *scaleR-nonneg*:
 assumes *c ≥ 0*
 shows *supermartingale M F t₀ (λ i ξ. c *_R X i ξ)*
proof
 {
 fix *i j :: 'b* **assume** *asm: t₀ ≤ i i ≤ j*
 thus *AE ξ in M. c *_R X i ξ ≥ cond-exp M (F i) (λ ξ. c *_R X j ξ) ξ*
 using *sigma-finite-subalgebra.cond-exp-scaleR-right*[*OF - integrable*, of *F i*
j c] *supermartingale-property*[*OF asm*] **by** (*fastforce intro!*: *scaleR-left-mono*[*OF -*

```

assms])
}
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)

lemma scaleR-le-zero:
  assumes  $c \leq 0$ 
  shows submartingale  $M F t_0 (\lambda i \xi. c *_{\mathbb{R}} X i \xi)$ 
proof
  {
    fix  $i j :: 'b$  assume asm:  $t_0 \leq i i \leq j$ 
    thus  $AE \xi \text{ in } M. c *_{\mathbb{R}} X i \xi \leq \text{cond-exp } M (F i) (\lambda \xi. c *_{\mathbb{R}} X j \xi) \xi$ 
      using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of  $F i j c$ ]
      supermartingale-property[OF asm] by (fastforce intro!: scaleR-left-mono-neg[OF -
      assms])
  }
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable
random-variable adapted borel-measurable-const-scaleR)

lemma uminus[intro]:
  shows submartingale  $M F t_0 (- X)$ 
  unfolding fun-Comp-def using scaleR-le-zero[of  $-1$ ] by simp

end

context supermartingale-linorder
begin

lemma set-integral-ge:
  assumes  $A \in F i t_0 \leq i i \leq j$ 
  shows set-lebesgue-integral  $M A (X i) \geq \text{set-lebesgue-integral } M A (X j)$ 
  using supermartingale-property[OF assms(2), of  $j$ ] assms subalgebras
  by (subst sigma-finite-subalgebra.cond-exp-set-integral[OF - integrable assms(1),
  of  $j$ ])
  (auto intro!: scaleR-left-mono integral-mono-AE-banach integrable-mult-indicator
  integrable simp add: subalgebra-def set-lebesgue-integral-def)

lemma min:
  assumes supermartingale-linorder  $M F t_0 Y$ 
  shows supermartingale-linorder  $M F t_0 (\lambda i \xi. \min (X i \xi) (Y i \xi))$ 
proof (unfold-locales)
  interpret  $Y: \text{supermartingale-linorder } M F t_0 Y$  by (rule assms)
  {
    fix  $i j :: 'b$  assume asm:  $t_0 \leq i i \leq j$ 
    have  $AE \xi \text{ in } M. \min (X i \xi) (Y i \xi) \geq \min (\text{cond-exp } M (F i) (X j) \xi) (\text{cond-exp } M (F i) (Y j) \xi)$ 
      using supermartingale-property  $Y.$ supermartingale-property asm
      unfolding min-def by fastforce
    thus  $AE \xi \text{ in } M. \min (X i \xi) (Y i \xi) \geq \text{cond-exp } M (F i) (\lambda \xi. \min (X j \xi) (Y j \xi)) \xi$ 
      using sigma-finite-subalgebra.cond-exp-min[OF - integrable  $Y.$ integrable, of

```

$F\ i\ j\ j]$ *asm* **by** (*fast intro: order.trans*)
}
show $\bigwedge i. t_0 \leq i \implies (\lambda \xi. \min (X\ i\ \xi) (Y\ i\ \xi)) \in \text{borel-measurable } (F\ i) \bigwedge i.$
 $t_0 \leq i \implies \text{integrable } M\ (\lambda \xi. \min (X\ i\ \xi) (Y\ i\ \xi))$ **by** (*force intro: Y.integrable*
integrable assms)
qed

lemma *min-0*:

shows *supermartingale-linorder* $M\ F\ t_0\ (\lambda i\ \xi. \min\ 0\ (X\ i\ \xi))$
proof –
interpret *zero: martingale-linorder* $M\ F\ t_0\ \lambda -.\ 0$ **by** (*force intro: martin-*
gale-linorder.intro)
show *?thesis* **by** (*intro zero.min supermartingale-linorder.intro supermartin-*
gale-axioms)
qed

end

lemma (*in sigma-finite-adapted-process-order*) *supermartingale-of-cond-exp-diff-le-zero*:

assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M\ (X\ i)$
and *diff-le-zero*: $\bigwedge i\ j. t_0 \leq i \implies i \leq j \implies \text{AE } x \text{ in } M. \text{ cond-exp } M\ (F\ i)$
 $(\lambda \xi. X\ j\ \xi - X\ i\ \xi)\ x \leq 0$
shows *supermartingale* $M\ F\ t_0\ X$
proof
{
fix $i\ j :: 'b$ **assume** *asm*: $t_0 \leq i \leq j$
thus *AE* ξ *in* $M. X\ i\ \xi \geq \text{cond-exp } M\ (F\ i)\ (X\ j)\ \xi$
using *diff-le-zero*[*OF asm*] *sigma-finite-subalgebra.cond-exp-diff*[*OF - inte-*
grable(1,1), of F i j i]
sigma-finite-subalgebra.cond-exp-F-meas[*OF - integrable adapted, of i*] **by**
fastforce
}
qed (*intro integrable*)

lemma (*in sigma-finite-adapted-process-linorder*) *supermartingale-of-set-integral-ge*:

assumes *integrable*: $\bigwedge i. t_0 \leq i \implies \text{integrable } M\ (X\ i)$
and $\bigwedge A\ i\ j. t_0 \leq i \implies i \leq j \implies A \in F\ i \implies \text{set-lebesgue-integral } M\ A\ (X\ j) \leq \text{set-lebesgue-integral } M\ A\ (X\ i)$
shows *supermartingale* $M\ F\ t_0\ X$
proof –
interpret *-: adapted-process* $M\ F\ t_0\ -X$ **by** (*rule uminus-adapted*)
interpret *uminus-X*: *sigma-finite-adapted-process-linorder* $M\ F\ t_0\ -X$..
note $*$ = *set-integral-uminus*[*unfolded set-integrable-def, OF integrable-mult-indicator*[*OF*
- integrable]]
have *supermartingale* $M\ F\ t_0\ (-(-X))$
using *ord-eq-le-trans*[*OF * ord-le-eq-trans*[*OF le-imp-neg-le*[*OF assms(2)*]] **[symmetric]*]]
subalgebras
by (*intro submartingale.uminus uminus-X.submartingale-of-set-integral-le*)
(clarsimp simp add: fun-Compl-def subalgebra-def integrable | fastforce)+

thus *?thesis* unfolding *fun-Compl-def* by *simp*
qed

7.8 Discrete Time Martingales

locale *nat-martingale* = *martingale* *M F 0 :: nat X* for *M F X*
locale *nat-submartingale* = *submartingale* *M F 0 :: nat X* for *M F X*
locale *nat-supermartingale* = *supermartingale* *M F 0 :: nat X* for *M F X*

locale *nat-submartingale-linorder* = *submartingale-linorder* *M F 0 :: nat X* for *M F X*
locale *nat-supermartingale-linorder* = *supermartingale-linorder* *M F 0 :: nat X* for *M F X*

sublocale *nat-submartingale-linorder* \subseteq *nat-submartingale* ..
sublocale *nat-supermartingale-linorder* \subseteq *nat-supermartingale* ..

lemma (in *nat-martingale*) *predictable-const*:
assumes *nat-predictable-process* *M F X*
shows *AE* ξ in *M*. *X i* ξ = *X j* ξ
proof –
have *: *AE* ξ in *M*. *X i* ξ = *X 0* ξ for *i*
proof (induction *i*)
case 0
then show ?case by (simp add: *bot-nat-def*)
next
case (Suc *i*)
interpret *S*: *nat-adapted-process* *M F* λi . *X (Suc i)* by (intro *nat-predictable-process.adapted-Suc* *assms*)
show ?case using *Suc S.adapted*[of *i*] *martingale-property*[*OF - le-SucI*, of *i*]
sigma-finite-subalgebra.cond-exp-F-meas[*OF - integrable*, of *F i Suc i*] by *fastforce*
qed
show ?thesis using *[of *i*] *[of *j*] by *force*
qed

lemma (in *nat-sigma-finite-adapted-process*) *martingale-of-set-integral-eq-Suc*:
assumes *integrable*: $\bigwedge i$. *integrable* *M (X i)*
and $\bigwedge A$ *i*. *A* $\in F i \implies$ *set-lebesgue-integral* *M A (X i)* = *set-lebesgue-integral* *M A (X (Suc i))*
shows *nat-martingale* *M F X*
proof (intro *nat-martingale.intro* *martingale-of-set-integral-eq*)
fix *i j A* assume *asm*: *i* $\leq j$ *A* \in *sets (F i)*
show *set-lebesgue-integral* *M A (X i)* = *set-lebesgue-integral* *M A (X j)* using
asm
proof (induction *j - i* arbitrary: *i j*)
case 0
then show ?case using *asm* by *simp*
next
case (Suc *n*)

hence *: $n = j - \text{Suc } i$ **by** *linarith*
 have $\text{Suc } i \leq j$ **using** *Suc(2,3)* **by** *linarith*
 thus ?case **using** *sets-F-mono[OF - le-SucI]* *Suc(4)* *Suc(1)[OF *]* **by** (*auto*
intro: assms(2)[THEN trans])
 qed
 qed (*simp add: integrable*)

lemma (*in nat-sigma-finite-adapted-process*) *martingale-nat*:
 assumes *integrable*: $\bigwedge i. \text{integrable } M (X i)$
 and $\bigwedge i. AE \xi \text{ in } M. X i \xi = \text{cond-exp } M (F i) (X (\text{Suc } i)) \xi$
 shows *nat-martingale* $M F X$
proof (*unfold-locales*)
 fix $i j :: \text{nat}$ **assume** *asm*: $i \leq j$
 show $AE \xi \text{ in } M. X i \xi = \text{cond-exp } M (F i) (X j) \xi$ **using** *asm*
proof (*induction j - i arbitrary: i j*)
 case 0
 hence $j = i$ **by** *simp*
 thus ?case **using** *sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted,*
THEN AE-symmetric] **by** *blast*
 next
 case (*Suc n*)
 have $j: j = \text{Suc } (n + i)$ **using** *Suc* **by** *linarith*
 have $n: n = n + i - i$ **using** *Suc* **by** *linarith*
 have *: $AE \xi \text{ in } M. \text{cond-exp } M (F (n + i)) (X j) \xi = X (n + i) \xi$ **unfolding**
j **using** *assms(2)[THEN AE-symmetric]* **by** *blast*
 have $AE \xi \text{ in } M. \text{cond-exp } M (F i) (X j) \xi = \text{cond-exp } M (F i) (\text{cond-exp } M$
 $(F (n + i)) (X j)) \xi$ **by** (*intro cond-exp-nested-subalg integrable subalg, simp add:*
subalgebra-def sets-F-mono)
 hence $AE \xi \text{ in } M. \text{cond-exp } M (F i) (X j) \xi = \text{cond-exp } M (F i) (X (n + i))$
 ξ **using** *cond-exp-cong-AE[OF integrable-cond-exp integrable *]* **by** *force*
 thus ?case **using** *Suc(1)[OF n]* **by** *fastforce*
 qed
 qed (*simp add: integrable*)

lemma (*in nat-sigma-finite-adapted-process*) *martingale-of-cond-exp-diff-Suc-eq-zero*:
 assumes *integrable*: $\bigwedge i. \text{integrable } M (X i)$
 and $\bigwedge i. AE \xi \text{ in } M. \text{cond-exp } M (F i) (\lambda \xi. X (\text{Suc } i) \xi - X i \xi) \xi = 0$
 shows *nat-martingale* $M F X$
proof (*intro martingale-nat integrable*)
 fix i
 show $AE \xi \text{ in } M. X i \xi = \text{cond-exp } M (F i) (X (\text{Suc } i)) \xi$ **using** *cond-exp-diff[OF*
integrable(1,1), of i Suc i i] *cond-exp-F-meas[OF integrable adapted, of i]* *assms(2)[of*
i] **by** *fastforce*
 qed

7.9 Discrete Time Submartingales

lemma (*in nat-submartingale*) *predictable-mono*:
 assumes *nat-predictable-process* $M F X i \leq j$

shows $AE \xi$ in M . $X i \xi \leq X j \xi$
using $assms(2)$
proof (*induction* $j - i$ arbitrary: $i j$)
 case 0
 then show ?*case* **by** *simp*
next
 case ($Suc\ n$)
 hence *: $n = j - Suc\ i$ **by** *linarith*
 interpret S : *nat-adapted-process* $M\ F\ \lambda i. X\ (Suc\ i)$ **by** (*intro nat-predictable-process.adapted-Suc*
assms)
 have $Suc\ i \leq j$ **using** $Suc(2,3)$ **by** *linarith*
 thus ?*case* **using** $Suc(1)[OF\ *]$ $S.adapted[of\ i]$ *submartingale-property*[$OF -$
le-SucI, *of i*] *sigma-finite-subalgebra.cond-exp-F-meas*[$OF - integrable$, *of F i Suc*
i] **by** *fastforce*
qed

lemma (*in nat-sigma-finite-adapted-process-linorder*) *submartingale-of-set-integral-le-Suc*:
 assumes *integrable*: $\bigwedge i. integrable\ M\ (X\ i)$
 and $\bigwedge A\ i. A \in F\ i \implies set-lebesgue-integral\ M\ A\ (X\ i) \leq set-lebesgue-integral$
M A (X (Suc i))
 shows *nat-submartingale* $M\ F\ X$
proof (*intro nat-submartingale.intro submartingale-of-set-integral-le*)
 fix $i\ j\ A$ **assume** *asm*: $i \leq j\ A \in sets\ (F\ i)$
 show $set-lebesgue-integral\ M\ A\ (X\ i) \leq set-lebesgue-integral\ M\ A\ (X\ j)$ **using**
asm
 proof (*induction* $j - i$ arbitrary: $i j$)
 case 0
 then show ?*case* **using** *asm* **by** *simp*
 next
 case ($Suc\ n$)
 hence *: $n = j - Suc\ i$ **by** *linarith*
 have $Suc\ i \leq j$ **using** $Suc(2,3)$ **by** *linarith*
 thus ?*case* **using** *sets-F-mono*[$OF - le-SucI$] $Suc(4)$ $Suc(1)[OF\ *]$ **by** (*auto*
intro: assms(2)[THEN order-trans])
 qed
qed (*simp add: integrable*)

lemma (*in nat-sigma-finite-adapted-process-linorder*) *submartingale-nat*:
 assumes *integrable*: $\bigwedge i. integrable\ M\ (X\ i)$
 and $\bigwedge i. AE\ \xi$ in M . $X\ i\ \xi \leq cond-exp\ M\ (F\ i)\ (X\ (Suc\ i))\ \xi$
 shows *nat-submartingale* $M\ F\ X$
 using *subalg integrable assms(2)*
 by (*intro submartingale-of-set-integral-le-Suc ord-le-eq-trans*[$OF\ set-integral-mono-AE-banach$
cond-exp-set-integral[symmetric]], *simp*)
 (*meson in-mono integrable-mult-indicator set-integrable-def subalgebra-def, me-*
son integrable-cond-exp in-mono integrable-mult-indicator set-integrable-def subal-
gebra-def, fast+)

lemma (*in nat-sigma-finite-adapted-process-linorder*) *submartingale-of-cond-exp-diff-Suc-nonneg*:

assumes *integrable*: $\bigwedge i. \text{integrable } M (X i)$
and $\bigwedge i. AE \xi \text{ in } M. \text{cond-exp } M (F i) (\lambda \xi. X (Suc i) \xi - X i \xi) \xi \geq 0$
shows *nat-submartingale* $M F X$
proof (*intro submartingale-nat integrable*)
fix i
show $AE \xi \text{ in } M. X i \xi \leq \text{cond-exp } M (F i) (X (Suc i)) \xi$ **using** *cond-exp-diff*[*OF integrable(1,1), of i Suc i i*] *cond-exp-F-meas*[*OF integrable adapted, of i*] *assms(2)*[*of i*] **by** *fastforce*
qed

lemma (*in nat-submartingale-linorder*) *partial-sum-scaleR*:

assumes *nat-adapted-process* $M F C \bigwedge i. AE \xi \text{ in } M. 0 \leq C i \xi \bigwedge i. AE \xi \text{ in } M. C i \xi \leq R$
shows *nat-submartingale* $M F (\lambda n \xi. \sum i < n. C i \xi *_R (X (Suc i) \xi - X i \xi))$
proof –
interpret C : *nat-adapted-process* $M F C$ **by** (*rule assms*)
interpret C' : *nat-adapted-process* $M F \lambda i \xi. C (i - 1) \xi *_R (X i \xi - X (i - 1) \xi)$ **by** (*intro nat-adapted-process.intro adapted-process.scaleR-right-adapted adapted-process.diff-adapted, unfold-locales*) (*auto intro: adaptedD C.adaptedD*) +
interpret C'' : *nat-adapted-process* $M F \lambda n \xi. \sum i < n. C i \xi *_R (X (Suc i) \xi - X i \xi)$ **by** (*rule C'.partial-sum-Suc-adapted[unfolded diff-Suc-1]*)
interpret S : *nat-sigma-finite-adapted-process-linorder* $M F (\lambda n \xi. \sum i < n. C i \xi *_R (X (Suc i) \xi - X i \xi))$..
have *integrable* $M (\lambda x. C i x *_R (X (Suc i) x - X i x))$ **for** i **using** *assms(2,3)*[*of i*] **by** (*intro Bochner-Integration.integrable-bound*[*OF integrable-scaleR-right, OF Bochner-Integration.integrable-diff, OF integrable(1,1), of R Suc i i*] (*auto simp add: mult-mono*))
moreover have $AE \xi \text{ in } M. 0 \leq \text{cond-exp } M (F i) (\lambda \xi. (\sum i < Suc i. C i \xi *_R (X (Suc i) \xi - X i \xi)) - (\sum i < i. C i \xi *_R (X (Suc i) \xi - X i \xi))) \xi$ **for** i
using *sigma-finite-subalgebra.cond-exp-measurable-scaleR*[*OF - calculation - C.adapted, of i*]
cond-exp-diff-nonneg[*OF - le-SucI, OF - order.refl, of i*] *assms(2,3)*[*of i*]
by (*fastforce simp add: scaleR-nonneg-nonneg integrable*)
ultimately show *?thesis* **by** (*intro S.submartingale-of-cond-exp-diff-Suc-nonneg Bochner-Integration.integrable-sum, blast+*)
qed

lemma (*in nat-submartingale-linorder*) *partial-sum-scaleR'*:

assumes *nat-predictable-process* $M F C \bigwedge i. AE \xi \text{ in } M. 0 \leq C i \xi \bigwedge i. AE \xi \text{ in } M. C i \xi \leq R$
shows *nat-submartingale* $M F (\lambda n \xi. \sum i < n. C (Suc i) \xi *_R (X (Suc i) \xi - X i \xi))$
proof –
interpret C : *nat-predictable-process* $M F C$ **by** (*rule assms*)
interpret $Suc-C$: *nat-adapted-process* $M F \lambda i. C (Suc i)$ **using** $C.adapted-Suc$.
show *?thesis* **by** (*intro partial-sum-scaleR*[*of - R*] *assms*) (*intro-locales*)
qed

7.10 Discrete Time Supermartingales

lemma (in *nat-supermartingale*) *predictable-mono*:

assumes *nat-predictable-process* $M \ F \ X \ i \leq j$

shows $AE \ \xi \text{ in } M. \ X \ i \ \xi \geq X \ j \ \xi$

using *assms*(2)

proof (*induction* $j - i$ arbitrary: $i \ j$)

case 0

then show ?*case* **by** *simp*

next

case (*Suc* n)

hence *: $n = j - \text{Suc } i$ **by** *linarith*

interpret S : *nat-adapted-process* $M \ F \ \lambda i. \ X \ (\text{Suc } i)$ **by** (*intro nat-predictable-process.adapted-Suc assms*)

have $\text{Suc } i \leq j$ **using** *Suc*(2,3) **by** *linarith*

thus ?*case* **using** *Suc*(1)[*OF* *] *S.adapted*[*of* i] *supermartingale-property*[*OF* - *le-SucI*, *of* i] *sigma-finite-subalgebra.cond-exp-F-meas*[*OF* - *integrable*, *of* $F \ i \ \text{Suc } i$] **by** *fastforce*

qed

lemma (in *nat-sigma-finite-adapted-process-linorder*) *supermartingale-of-set-integral-ge-Suc*:

assumes *integrable*: $\bigwedge i. \text{integrable } M \ (X \ i)$

and $\bigwedge A \ i. A \in F \ i \implies \text{set-lebesgue-integral } M \ A \ (X \ i) \geq \text{set-lebesgue-integral } M \ A \ (X \ (\text{Suc } i))$

shows *nat-supermartingale* $M \ F \ X$

proof -

interpret -: *adapted-process* $M \ F \ 0 \ -X$ **by** (*rule uminus-adapted*)

interpret *uminus-X*: *nat-sigma-finite-adapted-process-linorder* $M \ F \ -X$..

note * = *set-integral-uminus*[*unfolded set-integrable-def*, *OF integrable-mult-indicator*[*OF* - *integrable*]]

have *nat-supermartingale* $M \ F \ (-(-X))$

using *ord-eq-le-trans*[*OF* * *ord-le-eq-trans*[*OF le-imp-neg-le*[*OF assms*(2)] **[symmetric]*]] *subalgebras*

by (*intro nat-supermartingale.intro submartingale.uminus nat-submartingale.axioms uminus-X.submartingale-of-set-integral-le-Suc*)

(*clarsimp simp add: fun-Compl-def subalgebra-def integrable* | *fastforce*)+

thus ?*thesis* **unfolding** *fun-Compl-def* **by** *simp*

qed

lemma (in *nat-sigma-finite-adapted-process-linorder*) *supermartingale-nat*:

assumes *integrable*: $\bigwedge i. \text{integrable } M \ (X \ i)$

and $\bigwedge i. AE \ \xi \text{ in } M. \ X \ i \ \xi \geq \text{cond-exp } M \ (F \ i) \ (X \ (\text{Suc } i)) \ \xi$

shows *nat-supermartingale* $M \ F \ X$

proof -

interpret -: *adapted-process* $M \ F \ 0 \ -X$ **by** (*rule uminus-adapted*)

interpret *uminus-X*: *nat-sigma-finite-adapted-process-linorder* $M \ F \ -X$..

have $AE \ \xi \text{ in } M. \ -X \ i \ \xi \leq \text{cond-exp } M \ (F \ i) \ (\lambda x. -X \ (\text{Suc } i) \ x) \ \xi$ **for** i **using** *assms*(2) *cond-exp-uminus*[*OF integrable*, *of* $i \ \text{Suc } i$] **by** *force*

hence *nat-supermartingale* $M \ F \ (-(-X))$ **by** (*intro nat-supermartingale.intro submartingale.uminus nat-submartingale.axioms uminus-X.submartingale-nat*) (*auto*)

simp add: fun-Compl-def integrable)
thus *?thesis* **unfolding** *fun-Compl-def* **by** *simp*
qed

lemma (*in nat-sigma-finite-adapted-process-linorder*) *supermartingale-of-cond-exp-diff-Suc-le-zero*:
assumes *integrable*: $\bigwedge i. \text{integrable } M \ (X \ i)$
and $\bigwedge i. AE \ \xi \text{ in } M. \text{cond-exp } M \ (F \ i) \ (\lambda \xi. X \ (Suc \ i) \ \xi - X \ i \ \xi) \ \xi \leq 0$
shows *nat-supermartingale* *M F X*
proof (*intro supermartingale-nat integrable*)
fix *i*
show $AE \ \xi \text{ in } M. X \ i \ \xi \geq \text{cond-exp } M \ (F \ i) \ (X \ (Suc \ i)) \ \xi$ **using** *cond-exp-diff[OF integrable(1,1), of i Suc i i] cond-exp-F-meas[OF integrable adapted, of i] assms(2)[of i]* **by** *fastforce*
qed
end

theory *Example-Coin-Toss*
imports *Martingale HOL-Probability.Stream-Space HOL-Probability.Probability-Mass-Function*
begin

We consider a probability space consisting of infinite sequences of coin tosses.

definition *bernoulli-stream* :: *real* \Rightarrow (*bool stream*) *measure* **where**
bernoulli-stream *p* = *stream-space* (*measure-pmf* (*bernoulli-pmf* *p*))

lemma *space-bernoulli-stream[simp]*: *space* (*bernoulli-stream* *p*) = *UNIV* **by** (*simp add: bernoulli-stream-def space-stream-space*)

We define the fortune of the player at time *n* to be the number of heads minus number of tails.

definition *fortune* :: *nat* \Rightarrow *bool stream* \Rightarrow *real* **where**
fortune *n* = ($\lambda s. \sum b \leftarrow \text{stake } (Suc \ n) \ s. \text{if } b \text{ then } 1 \text{ else } -1$)

definition *toss* :: *nat* \Rightarrow *bool stream* \Rightarrow *real* **where**
toss *n* = ($\lambda s. \text{if } snth \ s \ n \text{ then } 1 \text{ else } -1$)

lemma *toss-indicator-def*: *toss* *n* = *indicator* $\{s. s \ !! \ n\}$ - *indicator* $\{s. \neg s \ !! \ n\}$
unfolding *toss-def indicator-def* **by** *force*

lemma *range-toss*: *range* (*toss* *n*) = $\{-1, 1\}$

proof -
have *sconst True !! n* **by** *simp*
moreover **have** $\neg sconst \ False \ !! \ n$ **by** *simp*
ultimately **have** $\exists x. x \ !! \ n \ \exists x. \neg x \ !! \ n$ **by** *blast+*
thus *?thesis* **unfolding** *toss-def image-def* **by** *auto*
qed

lemma *vimage-toss*: $\text{toss } n - 'A = (\text{if } 1 \in A \text{ then } \{s. s !! n\} \text{ else } \{\}) \cup (\text{if } -1 \in A \text{ then } \{s. \neg s !! n\} \text{ else } \{\})$

unfolding *vimage-def toss-def* **by** *auto*

lemma *fortune-Suc*: $\text{fortune } (\text{Suc } n) s = \text{fortune } n s + \text{toss } (\text{Suc } n) s$

by (*induction n arbitrary: s*) (*simp add: fortune-def toss-def*)+

lemma *fortune-toss-sum*: $\text{fortune } n s = (\sum i \in \{..n\}. \text{toss } i s)$

by (*induction n arbitrary: s*) (*simp add: fortune-def toss-def, simp add: fortune-Suc*)

lemma *fortune-bound*: $\text{norm } (\text{fortune } n s) \leq \text{Suc } n$ **by** (*induction n*) (*force simp add: fortune-toss-sum toss-def*)+

Our definition of *bernoulli-stream* constitutes a probability space.

interpretation *prob-space bernoulli-stream p* **unfolding** *bernoulli-stream-def* **by** (*simp add: measure-pmf.prob-space-axioms prob-space.prob-space-stream-space*)

abbreviation *toss-filtration p* $\equiv \text{nat-natural-filtration } (\text{bernoulli-stream } p) \text{ toss}$

The stochastic process *toss* is adapted to the filtration it generates.

interpretation *toss: nat-adapted-process bernoulli-stream p nat-natural-filtration* (*bernoulli-stream p*) *toss toss*

by (*intro nat-adapted-process.intro stochastic-process.adapted-process-natural-filtration*) (*unfold-locales, auto simp add: toss-def bernoulli-stream-def*)

Similarly, the stochastic process *fortune* is adapted to the filtration generated by the tosses.

interpretation *fortune: nat-finite-adapted-process-linorder bernoulli-stream p nat-natural-filtration* (*bernoulli-stream p*) *toss fortune*

proof –

show *nat-finite-adapted-process-linorder* (*bernoulli-stream p*) (*toss-filtration p*) *fortune*

unfolding *fortune-toss-sum*

by (*intro nat-finite-adapted-process-linorder.intro*

finite-adapted-process-linorder.intro

finite-adapted-process-order.intro

finite-adapted-process.intro

toss.partial-sum-adapted[folded atMost-atLeast0]) intro-locales

qed

lemma *integrable-toss*: *integrable* (*bernoulli-stream p*) (*toss n*)

using *toss.random-variable*

by (*intro Bochner-Integration.integrable-bound[OF integrable-const[of - 1 :: real]]*) (*auto simp add: toss-def*)

lemma *integrable-fortune*: *integrable* (*bernoulli-stream p*) (*fortune n*) **using** *fortune-bound*

by (intro Bochner-Integration.integrable-bound[OF integrable-const[of - Suc n] fortune.random-variable]) auto

We provide the following lemma to explicitly calculate the probability of events in this probability space.

lemma *measure-bernoulli-stream-snth-pred*:

assumes $0 \leq p$ **and** $p \leq 1$ **and** finite J

shows $\text{prob } p \{w \in \text{space } (\text{bernoulli-stream } p). \forall j \in J. P j = w !! j\} = p^{\wedge \text{card } (J \cap \text{Collect } P)} * (1 - p)^{\wedge (\text{card } (J - \text{Collect } P))}$

proof –

let $?PiE = (\Pi_E i \in J. \text{if } P i \text{ then } \{\text{True}\} \text{ else } \{\text{False}\})$

have $\text{product-prob-space } (\lambda-. \text{measure-pmf } (\text{bernoulli-pmf } p))$ **by** *unfold-locales*

hence $*$: $\text{to-stream } -' \{s. \forall i \in J. P i = s !! i\} = \{s. \forall i \in J. P i = s i\}$ **using** *assms by (simp add: to-stream-def)*

also have $\dots = \text{prod-emb UNIV } (\lambda-. \text{measure-pmf } (\text{bernoulli-pmf } p)) J ?PiE$

proof –

{

fix s **assume** $(\forall i \in J. P i = s i)$

hence $(\forall i \in J. P i = s i) = (s \in \text{prod-emb UNIV } (\lambda-. \text{measure-pmf } (\text{bernoulli-pmf } p)) J ?PiE)$

by (*subst prod-emb-iff[of s] (smt (verit, best) not-def assms(3) id-def PiE-eq-singleton UNIV-I extensional-UNIV insert-iff singletonD space-measure-pmf)*)

}

moreover

{

fix s **assume** $\neg(\forall i \in J. P i = s i)$

then obtain i **where** $i \in J$ $P i \neq s i$ **by** *blast*

hence $(\forall i \in J. P i = s i) = (s \in \text{prod-emb UNIV } (\lambda-. \text{measure-pmf } (\text{bernoulli-pmf } p)) J ?PiE)$

by (*simp add: restrict-def prod-emb-iff[of s] (smt (verit, ccfv-SIG) PiE-mem assms(3) id-def insert-iff singleton-iff)*)

}

ultimately show *?thesis* **by** *auto*

qed

finally have *inteq*: $(\text{to-stream } -' \{s. \forall i \in J. P i = s !! i\}) = \text{prod-emb UNIV } (\lambda-. \text{measure-pmf } (\text{bernoulli-pmf } p)) J ?PiE$.

let $?M = (\Pi_M \text{ UNIV } (\lambda-. \text{measure-pmf } (\text{bernoulli-pmf } p)))$

have $\text{emeasure } (\text{bernoulli-stream } p) \{s \in \text{space } (\text{bernoulli-stream } p). \forall i \in J. P i = s !! i\} = \text{emeasure } ?M (\text{to-stream } -' \{s. \forall i \in J. P i = s !! i\})$

using *assms emeasure-distr[of to-stream ?M (vimage-algebra (streams (space (measure-pmf (bernoulli-pmf p)))) (!) ?M) {s. \forall i \in J. P i = s !! i}, symmetric] measurable-to-stream[of (measure-pmf (bernoulli-pmf p))]*

by (*simp only: bernoulli-stream-def stream-space-def *, simp add: space-PiM*) (*smt (verit, best) emeasure-notin-sets in-vimage-algebra inf-top.right-neutral sets-distr vimage-Collect*)

also have $\dots = \text{emeasure } ?M (\text{prod-emb UNIV } (\lambda-. \text{measure-pmf } (\text{bernoulli-pmf } p)) J ?PiE)$ **using** *inteq* **by** (*simp add: space-PiM*)

also have $\dots = (\prod_{i \in J. \text{emeasure } (\text{measure-pmf } (\text{bernoulli-pmf } p))} (\text{if } P i \text{ then$

```

{ True } else { False }))
  by (subst emeasure-PiM-emb) (auto simp add: prob-space-measure-pmf assms(3))
  also have ... = ( $\prod_{i \in J} \text{Collect } P. \text{ennreal } p$ ) * ( $\prod_{i \in J - \text{Collect } P} \text{ennreal } (1 - p)$ )
  unfolding emeasure-pmf-single[of bernoulli-pmf p True, unfolded pmf-bernoulli-True[OF assms(1,2)], symmetric]
    emeasure-pmf-single[of bernoulli-pmf p False, unfolded pmf-bernoulli-False[OF assms(1,2)], symmetric]
  by (simp add: prod.Int-Diff[OF assms(3), of - Collect P])
  also have ... =  $p^{\text{card } (J \cap \text{Collect } P)} * (1 - p)^{\text{card } (J - \text{Collect } P)}$  using
    assms by (simp add: prod-ennreal ennreal-mult' ennreal-power)
  finally show ?thesis using assms by (intro measure-eq-emeasure-eq-ennreal)
auto
qed

```

```

lemma
  assumes  $0 \leq p$  and  $p \leq 1$ 
  shows measure-bernoulli-stream-snth:  $\text{prob } p \{w \in \text{space } (\text{bernoulli-stream } p). w !! i\} = p$ 
  and measure-bernoulli-stream-neg-snth:  $\text{prob } p \{w \in \text{space } (\text{bernoulli-stream } p). \neg w !! i\} = 1 - p$ 
  using measure-bernoulli-stream-snth-pred[OF assms, of  $\{i\} \lambda x. \text{True}$ ]
    measure-bernoulli-stream-snth-pred[OF assms, of  $\{i\} \lambda x. \text{False}$ ] by auto

```

Now we can express the expected value of a single coin toss.

```

lemma integral-toss:
  assumes  $0 \leq p$  and  $p \leq 1$ 
  shows expectation p (toss n) =  $2 * p - 1$ 
proof -
  have [simp]:  $\{s. s !! n\} \in \text{events } p$  using measurable-snth[THEN measurable-sets,
    of {True} measure-pmf (bernoulli-pmf p) n, folded bernoulli-stream-def]
  by (simp add: vimage-def)
  have expectation p (toss n) = Bochner-Integration.simple-bochner-integral (bernoulli-stream
    p) (toss n)
  using toss.random-variable[of n, THEN measurable-sets]
  by (intro simple-bochner-integrable-eq-integral[symmetric] simple-bochner-integrable.intros)
    (auto simp add: toss-def simple-function-def image-def)
  also have ... =  $p - \text{prob } p \{s. \neg s !! n\}$  unfolding simple-bochner-integral-def
  using measure-bernoulli-stream-snth[OF assms]
  by (simp add: range-toss, simp add: toss-def)
  also have ... =  $p - (1 - \text{prob } p \{s. s !! n\})$  by (subst prob-compl[symmetric],
    auto simp add: Collect-neg-eq Compl-eq-Diff-UNIV)
  finally show ?thesis using measure-bernoulli-stream-snth[OF assms] by simp
qed

```

Now, we show that the tosses are independent from one another.

```

lemma indep-vars-toss:
  assumes  $0 \leq p$  and  $p \leq 1$ 
  shows indep-vars p ( $\lambda \cdot. \text{borel}$ ) toss  $\{0..\}$ 

```


proof (*subst indep-vars-def, intro conjI indep-sets-sigma*)
 {
 fix $A\ J$ **assume** $asm: J \neq \{\}$ *finite* $J\ \forall j \in J. A\ j \in \{\text{toss } j - 'A \cap \text{space}$
(bernoulli-stream p) | A. A ∈ borel}
 hence $\forall j \in J. \exists B \in \text{borel}. A\ j = \text{toss } j - 'B \cap \text{space (bernoulli-stream p)}$ **by**
auto
 then obtain B **where** $B\text{-is}: A\ j = \text{toss } j - 'B\ j \cap \text{space (bernoulli-stream p)}$
B j ∈ borel if j ∈ J for j by metis

 have $\text{prob } p (\bigcap (A - 'J)) = (\prod_{j \in J}. \text{prob } p (A\ j))$
proof *cases*

We consider the case where there is a zero probability event.

assume $\exists j \in J. 1 \notin B\ j \wedge -1 \notin B\ j$
 then obtain j **where** $j\text{-is}: j \in J\ 1 \notin B\ j\ -1 \notin B\ j$ **by** *blast*
hence $A\text{-j-empty}: A\ j = \{\}$ **using** $B\text{-is}$ **by** *(force simp add: toss-def vimage-def)*
 hence $\bigcap (A - 'J) = \{\}$ **using** $j\text{-is}$ **by** *blast*
 moreover have $\text{prob } p (A\ j) = 0$ **using** $A\text{-j-empty}$ **by** *simp*
 ultimately show *?thesis* **using** $j\text{-is}$ $asm(2)$ **by** *auto*
next

We now assume all events have positive probability.

assume $\neg(\exists j \in J. 1 \notin B\ j \wedge -1 \notin B\ j)$
 hence $*$: $1 \in B\ j \vee -1 \in B\ j$ **if** $j \in J$ **for** j **using** *that* **by** *blast*

define J' **where** $[simp]: J' = \{j \in J. (1 \in B\ j) \longleftrightarrow (-1 \notin B\ j)\}$
 hence $\text{toss } j\ w \in B\ j \longleftrightarrow (1 \in B\ j) = w !! j$ **if** $j \in J'$ **for** $w\ j$ **using** *that*
unfolding *toss-def* **by** *simp*
 hence $(\bigcap (A - 'J')) = \{w \in \text{space (bernoulli-stream p)}. \forall j \in J'. (1 \in B\ j) =$
w !! j} using B-is by force
 hence $\text{prob-}J': \text{prob } p (\bigcap (A - 'J')) = p \wedge \text{card } (J' \cap \{j. 1 \in B\ j\}) * (1 -$
p) ^ card (J' - {j. 1 ∈ B j})
 using *measure-bernoulli-stream-snth-pred[OF assms finite-subset[OF -*
asm(2)], of J' λj. 1 ∈ B j] **by** *auto*

The index set J' consists of the indices of all non-trivial events.

have $A\text{-j-True}: A\ j = \{w \in \text{space (bernoulli-stream p)}. w !! j\}$ **if** $j \in J' \cap \{j.$
1 ∈ B j} for j
 using *that* **by** *(auto simp add: toss-def B-is(1) split: if-splits)*

have $A\text{-j-False}: A\ j = \{w \in \text{space (bernoulli-stream p)}. \neg w !! j\}$ **if** $j \in J' -$
{j. 1 ∈ B j} for j
 using *that* $B\text{-is}$ **by** *(auto simp add: toss-def)*

have $A\text{-j-top}: A\ j = \text{space (bernoulli-stream p)}$ **if** $j \in J - J'$ **for** j **using** *that*
** by (auto simp add: B-is toss-def)*
 hence $\bigcap (A - 'J) = \bigcap (A - 'J')$ **by** *auto*
 hence $\text{prob } p (\bigcap (A - 'J)) = \text{prob } p (\bigcap (A - 'J'))$ **by** *presburger*

also have ... = $(\prod_{j \in J' \cap \{j. 1 \in B\}} \text{prob } p(A\ j)) * (\prod_{j \in J' - \{j. 1 \in B\}} \text{prob } p(A\ j))$
by (*simp only: prob-J' A-j-True A-j-False measure-bernoulli-stream-snth*[*OF assms*] *measure-bernoulli-stream-neg-snth*[*OF assms*] *cong: prod.cong*) *simp*
also have ... = $(\prod_{j \in J'} \text{prob } p(A\ j))$ **using** *asm(2)* **by** (*intro prod.Int-Diff*[*symmetric*]) *auto*
also have ... = $(\prod_{j \in J'} \text{prob } p(A\ j)) * (\prod_{j \in J - J'} \text{prob } p(A\ j))$ **using** *A-j-top prob-space* **by** *simp*
also have ... = $(\prod_{j \in J} \text{prob } p(A\ j))$ **using** *asm(2)* **by** (*metis* (*no-types*, *lifting*) *J'-def mem-Collect-eq mult commute prod.subset-diff subsetI*)
finally show ?thesis .
qed
}
thus *indep-sets* p ($\lambda i. \{ \text{toss } i - 'A \cap \text{space } (\text{bernoulli-stream } p) \mid A. A \in \text{sets borel} \} \{0..\}$) **using** *measurable-sets*[*OF toss.random-variable*]
by (*intro indep-setsI subsetI*) *fastforce*
qed (*simp*, *intro Int-stableI*, *simp*, *metis sets.Int vimage-Int*)

The fortune of a player is a martingale (resp. sub- or supermartingale) with respect to the filtration generated by the coin tosses.

theorem *fortune-martingale*:

assumes $p = 1/2$
shows *nat-martingale* (*bernoulli-stream* p) (*toss-filtration* p) *fortune*
using *cond-exp-indep*[*OF fortune.subalg indep-set-natural-filtration integrable-toss*, *OF zero-order(1) lessI indep-vars-toss*, *of p*]
integral-toss assms
by (*intro fortune.martingale-of-cond-exp-diff-Suc-eq-zero integrable-fortune*)
(force simp add: fortune-toss-sum)

theorem *fortune-submartingale*:

assumes $1/2 \leq p \leq 1$
shows *nat-submartingale* (*bernoulli-stream* p) (*toss-filtration* p) *fortune*
proof (*intro fortune.submartingale-of-cond-exp-diff-Suc-nonneg integrable-fortune*)
fix n
show $\text{AE } \xi \text{ in } \text{bernoulli-stream } p. 0 \leq \text{cond-exp } (\text{bernoulli-stream } p) (\text{toss-filtration } p\ n) (\lambda \xi. \text{fortune } (\text{Suc } n) \ \xi - \text{fortune } n \ \xi)$
using *cond-exp-indep*[*OF fortune.subalg indep-set-natural-filtration integrable-toss*, *OF zero-order(1) lessI indep-vars-toss*, *of p n*]
integral-toss[*of p Suc n*] *assms*
by (*force simp add: fortune-toss-sum*)
qed

theorem *fortune-supermartingale*:

assumes $0 \leq p \leq 1/2$
shows *nat-supermartingale* (*bernoulli-stream* p) (*toss-filtration* p) *fortune*
proof (*intro fortune.supermartingale-of-cond-exp-diff-Suc-le-zero integrable-fortune*)
fix n
show $\text{AE } \xi \text{ in } \text{bernoulli-stream } p. 0 \geq \text{cond-exp } (\text{bernoulli-stream } p) (\text{toss-filtration } p\ n) (\lambda \xi. \text{fortune } (\text{Suc } n) \ \xi - \text{fortune } n \ \xi)$

```

using cond-exp-indep[OF fortune.subalg indep-set-natural-filtration integrable-toss,
OF zero-order(1) lessI indep-vars-toss, of p n]
      integral-toss[of p Suc n] assms
by (force simp add: fortune-toss-sum)
qed

end

```