

Simulation of CPUID of Intel Architecture

Grigory Rechistov*

Name Surname[†]

March 20, 2014

Contents

1	Introduction	1
1.1	Contributions	2
2	Overview of Processor Identification	2
2.1	MIPS	2
2.2	ARM	3
2.3	IBM System z	3
2.4	PowerPC	3
2.5	SPARC	3
2.6	Intel IA-64 (Itanium)	4
2.7	Intel IA-32 and Intel 64	4
2.8	Comparison of Processor Identification	4
3	What is so hard about Intel CPUID	6
4	Existing Approaches to CPUID Simulation	7
4.1	Bochs	7
4.2	Xen	7
4.3	Qemu	7
4.4	Simics	7
5	CPUID Definition and Generation	7
6	Evaluation	7
6.1	Compatibility with existing code	7
6.2	Extensions	7
7	Conclusions	8

Abstract

In this paper we give an overview of existing microprocessor features identification facilities. Then we describe our approach to implementation of a software model of Intel IA-32 CPUID instruction. The described solution allows to define all recent CPUs' features, as well as future extensions. Our model was incorporated into the Wind River Simics^{*} simulator framework.

Key words: cpu identification, cpu simulation, CPUID, Simics.

1 Introduction

Most of documents defining CPU architectures do not specify every implementation detail; only interfaces, such as register state and instruction set (ISA), are defined. Microarchitectural details are usually left to vendors.

^{*}Moscow Institute of Physics and Technology, grigory.rechistov@phystech.edu

[†]email@mail.com

Any more-or-less mature processor architecture will have a set of extensions meant to improve performance, reliability, energy consumption or any other aspect of computer operation. A long evolution happened with backwards compatibility in mind and influenced by strong market competition will lead to a situation when there are quite a few of extensions, sometimes contradicting each other but nevertheless designated as “compatible” architectures.

Software programmers, who wish to use a certain architecture extension in their code, should be able to reliably identify its presence at run time. Therefore, processors provide means of feature identification in forms of instructions and/or registers which, being accessed, yield values that encode extended capabilities: vendor name, model name, machine word width, register count and similar model specific information.

Software simulators serve needs of pre-silicon software development, system firmware debugging, architecture and algorithm exploration etc. As with hardware, in order to be useful, software embodiments of processor specifications must faithfully declare supported extensions. Yet, building a software model for processor feature identification may turn out to be a non-trivial task, complicated by a bulk and complexity of information that needs to be represented, especially when a number of supported processor models is high.

From our experience with pre-silicon simulation tools, certain types of target software, such as BIOS, firmware and drivers, are quite sensitive to processor identification and will refuse to work or crash if a single bit of it is wrong. At the same time, such software in its early age may be incomplete; a software model has to adapt to mimic features that are absent for it to work. A naïve ad-hoc approach to feature data storage and handling will lead to a code that is very hard to extend support. These circumstances lead us to develop a dedicated framework for simulation of CPUID instruction, described in the paper.

The rest of this paper is organized as follows. In the section 2, we present a survey of approaches to defining model-specific features used in different microprocessor architectures. In the section 3 it is shown that the Intel IA-32 has one of the most complex CPUID instruction. Section 4 describes approaches to CPUID representation and simulation found in several software models. A new specification language for CPUID definition and tools for generation of simulation code are described in section 5. Our experience with its integration into an industrial simulation framework Wind River Simics* are shown in section 6. Section 7 gives final remarks and future work directions.

1.1 Contributions

In this paper we make the following contributions.

1. Comparison of existing approaches to extensions/feature identification of different modern processor architectures. To our knowledge, no such survey has been published yet.
2. Describe design, implementation and evaluation of a structured solution to the simulation of CPUID instruction of Intel IA-32. Analysis of several open source and proprietary simulators has shown that CPUID simulation is traditionally made ad-hoc, resulting in an entanglement of code which is hard to support.

2 Overview of Processor Identification

This section outlines mechanisms for processor identification employed by several modern general purpose microprocessor architectures. We do not attempt to give exhaustive descriptions of instruction sets; nor we intend to give a historical overview of now discontinued designs. The goal is to highlight differences and similarities between systems. Detailed explanations for all mentioned registers and instructions can be found at respective reference manuals cited below.

2.1 MIPS

MIPS CPUs store processor identification within PRid register, which is the 15th register in Coprocessor 0 [8]. It contains 32 bits of information (Fig. 1), part of which is vendor-specific.

31	24 23	16 15	8 7	0
Company Options		Company ID	Processor ID	Revision

Figure 1: MIPS PRid register fields

MIPS architecture also implements special register that contains information to list capabilities of floating point unit, termed Floating Point Implementation Register (FIR) [14]. It is a read-only 32-bit register (Fig. 2).

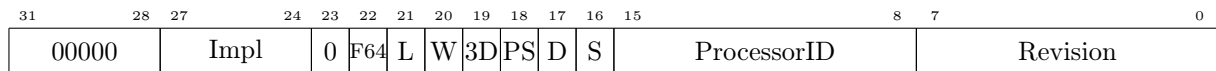


Figure 2: MIPS FIR register fields

TODO Check this
http://hwdb.mipt.cc/MIPS_PRIId_register
<http://code.google.com/p/phantomuserland/source/browse/trunk/phantom/dev/mips/cpuid.c?r=1094>
<http://www.imgtec.com/mips/mips32-architecture.asp>

2.2 ARM

The popular RISC architecture ARM provides two different ways to identified cores. This ways described in [2].

The first way is to read info form the Register 0 of the System Control Coprocessor also called CPUID base register [3]. As with MIPS, it holds 32 bits of information and look a bit differ for ARM7 core family (Fig. 3) and ARM9 and later core families (Fig. fig:arm-cpuid-v9).

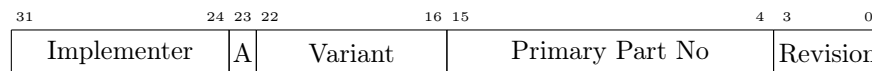


Figure 3: ARM7 Core Family CPUID base register fields

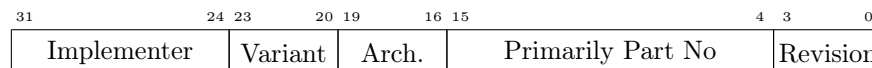


Figure 4: ARM9 and later Core Family CPUID base register fields

Examples of values [11]: Intel (XScale) PXA272 – 0x69054117, Qualcomm MSM7200A — 0x4117b362.

The other way that ARM core can be identified is through a TAP ID [2] that used ton configure debug software and indicates the information about an ARM core.

2.3 IBM System z

IBM System z10 [6] provides STSI and STIDP instructions. The information is rather scarce on the subject. **TODO** Expand

The STSI instruction used to determine processor model and model capacity identifier for the base configuration and for any additional changes through update actions.

The STIDP instruction provides information about the processor type, serial number, and logical partition identifier (Fig. 5).

2.4 PowerPC

PowerPC [7] offers a 32-bit PVR register that contains just version and revision numbers.

For ISA extensions a mechanism of APU (application processor units) is employed. MSR (machine state register) is used to store information of APUs available.

TODO See also
<http://lxr.linux.no/#linux+v3.13.5/arch/powerpc/kernel/cputable.c#L2245> — identify_cpu function
http://cache.freescale.com/files/archives/doc/support_info/PPCPVR.pdf

2.5 SPARC

The SPARC v9 standard [17] leaves quite a number of details implementation specific. To distinct between version a 64-bit register VER is defined (Fig. 6).

63	4847	3231	8 7	0
Logical partition 2-digit indicator	Machine type number	CPU identification number	Version code	

Figure 5: STIDP output

63	4847	3231	2423	1615	8 7	5 4	0
Manufacturer	Implementation	Revision	—	Max trap levels	—	Max window	

Figure 6: SPARC v9 VER register fields

2.6 Intel IA-64 (Itanium)

Intel IA-64, also known as Itanium™ [10] was conceived after the original Intel 80x86 series and was meant to supplant it. A set of CPUID registers is used for identification purposes. The set’s design somewhat resembles IA-32 CPUID instruction (discussed shortly after) — its register numbers loosely resemble leaves of the latter. At the moment of this writing, all announced IA-64 systems offered up to five CPUID registers. To have room for feature expansion, bits 0–7 of CPUID[3] store the actual size of the register set (limiting it to 256 entries).

A CPUID table for an Itanium 2 system is given on Fig. 7. Total capacity of it can be estimated as $5 \times 64 = 320$ bits.

2.7 Intel IA-32 and Intel 64

The common IBM PC architecture, starting from Intel Pentium™ and its clones, provides CPUID instruction [9, 1]. The 64-bit extension, known as Intel 64 or AMD64, is unchanged in this regard and we will make no distinction between them. CPUID takes two input operands in 32-bit registers EAX and ECX (called *leaf* and *subleaf*) and puts the result to four 32-bit registers, namely EAX, EBX, ECX and EDX.

Since its inception it has been extended numerous number of times. **TODO** Expand

On Figure 8 an output of all defined leaves and subleaves for a modern Intel processor (of microarchitecture Intel Ivy Bridge) is shown. The resulting table contains 25 tuples of 4 values of 32 bit width each, total of more than 3 kbit of information. Essentially a number of features can be deduced from it.

- Processor producer brand (leaf 0) and SKU naming (leaves 0x80000002–0x80000004).
- Availability of ISA extensions such as 64 bit mode, SSE2/3/4.1/4.2, AVX, MOVBE etc.
- Cache configuration of all layers, both in legacy format (leaf 2) and in current format (leaf 4 with subleaves).
- Multi-processor configuration knowledge, such as availability of Intel HyperThreading, relative position inside CPU package (topology at leaf 11), presence of APIC (interrupt controller).
- Numerous parameters for the implementation, such as addresses widths (leaf 0x80000008), availability of dynamic frequency scaling, supported debugging facilities etc.

2.8 Comparison of Processor Identification

Based on the presented data several conclusions can be made.

Leaf	Value
0	0x49656e69756e6547
0x1	0x6c65746e
0x2	0
0x3	0x20010104
0x4	0x5

Figure 7: Contents of CPUID registers for an Intel Itanium 9100 (code name Montvale) system using ggg-cpuid [16]

Leaf	Subleaf	EAX	EBX	ECX	EDX

0	0	0xd	0x756e6547	0x6c65746e	0x49656e69
0x1	0	0x306a9	0x6100800	0x7f9ae3bf	0xbfebfbff
0x2	0	0x76035a01	0xf0b0ff	0	0xca0000
0x4	0	0x1c004121	0x1c0003f	0x3f	0
0x4	0x1	0x1c004122	0x1c0003f	0x3f	0
0x4	0x2	0x1c004143	0x1c0003f	0x1ff	0
0x4	0x3	0x1c03c163	0x2c0003f	0x1fff	0x6
0x5	0	0x40	0x40	0x3	0x1120
0x6	0	0x77	0x2	0x9	0
0x7	0	0	0x281	0	0
0xa	0	0x7300803	0	0	0x603
0xb	0	0x1	0x1	0x100	0x6
0xb	0x1	0x4	0x4	0x201	0x6
0xb	0x2	0	0	0x2	0x6
0xb	0x3	0	0	0x3	0x6
0xd	0	0x7	0x340	0x340	0
0xd	0x1	0x1	0	0	0
0xd	0x2	0x100	0x240	0	0
0x80000000	0	0x80000008	0	0	0
0x80000001	0	0	0	0x1	0x28100800
0x80000002	0	0x20202020	0x20202020	0x65746e49	0x2952286c
0x80000003	0	0x726f4320	0x4d542865	0x35692029	0x3534332d
0x80000004	0	0x50432030	0x20402055	0x30312e33	0x7a4847
0x80000006	0	0	0	0x1006040	0
0x80000007	0	0	0	0	0x100
0x80000008	0	0x3024	0	0	0

Figure 8: Output of CPUID instruction obtained for Intel® Core™ i5-3450 using ggg-cpuid [16]

```

flags :  fpu vme de pse tsc msr pae mce cx8 apic sep mtrr pge mca cmov pat pse36 clflush
dts acpi mmx fxsr sse sse2 ss ht tm pbe syscall nx rdtscp lm constant_tsc arch_perfmon
pebs bts nopl xtopology nonstop_tsc aperfmperf pni pclmulqdq dtes64 monitor ds_cpl vmx
est tm2 ssse3 cx16 xtpr pdcm pcid sse4_1 sse4_2 x2apic popcnt tsc_deadline_timer xsave
avx lahf_lm arat epb xsaveopt pln pts dtherm tpr_shadow vnmi flexpriority ept vpid

```

Figure 9: Part of output of `cat /proc/cpuinfo` command on GNU/Linux on a recent Intel IA-32 CPU

- CPU identification facilities differ greatly between architectures. They may be represented by instructions, registers, or groups of both.
- The most common thing that can be specified through a CPUID is a vendor identification. The next on popularity is indication of ISA extensions.
- The complexity/completeness of CPUID facilities depends on whether there are requirements to run the same binary code on hardware from multiple vendors and/or of different generations. If, in order to perform efficiently, software must “know” a list of supported instruction extensions and other types of model specific configuration, there has to be a documented way to obtain such knowledge.
- Conversely, CPUs provided by a single vendor and/or designed for specific software usually provide less means of self-identification. Most microcontrollers for embedded applications hardly provide even an idea of CPUID, compared to general purpose processors, because software is often written to be run on a specific HW; binaries are not meant to be moved to some other incarnation of the same architecture.

3 What is so hard about Intel CPUID

We now concentrate solely on the Intel IA-32 architecture and the single CPUID instruction. A complete definition of CPUID in [9] takes about 40 pages.

There is a number of complications that have resulted from long uncontrolled expansion of the CPUID.

Influence of multiple vendors. Until recently, there were numerous CPU vendors that offered processors compatible with Intel architecture, including IBM, Cyrix, VIA, Centaur, Transmeta etc. By 2014, considerably fewer companies remain — in particular Intel itself and AMD. A coordination committee to define how new IA-32 extensions are indicated have not existed. Uncontrolled competition lead to the numerous small and subtle but essential differences in implementation. Until (and for some time after) CPUID was introduced, a robust identification of IA-32 processor brand and model required surprisingly intricate methods [5]. At present times, things are somewhat better documented, but still are not controlled in a centralized manner.

Long history that lead to extremely long list of extensions. With 40 years of backwards-compatible development the IA-32 architecture collected a number of additions. Consider a list of flags which a modern GNU/Linux operating system shows for a CPUID of the laptop this paper is being written on (Fig. 9). It should be noted that only a part of information from CPUID is actually present on this list.

Now, moving to details of CPUID operation, we will show more complexity that arises from the semantics of CPUID instruction.

Elements addressing To inspect a value of a particular

- Leaves
- Subleaves
- Registers
- Bit range

Non-constant values Firmware is able to suppress certain features indicated by CPUID by manipulating bits of model specific register (MSR) IA32_MISC_ENABLE. CPUID leaves more than 0x3 and less than 0x80000000 are visible only when IA32_MISC_ENABLE.BOOT_NT4[bit 22] is clear. For example, operating system of BIOS may disable MWAIT instruction by using IA32_MISC_ENABLE MSR; disabling MWAIT also clears corresponding CPUID feature flag. Software is also able to manage several architecture extensions and CPUID flags using control register CR4. For example operating system can set OSXSAVE flag to indicate that use of XGETBV, XSAVE and XRESTOR instructions is supported by general software.

Topology-variable elements Finally, it should be noted that, besides EAX, EBX, ECX, EDX, one more register may be affected by CPUID, namely IA32_SIGNATURE **TODO**.

4 Existing Approaches to CPUID Simulation

Every software simulator, emulator or virtual machine of a recent (Pentium or later) IA-32 system must guarantee CPUID operation of certain accuracy. If such tool is used for system firmware development, which is even more sensitive to identification information, the following requirements should be.

- Be accurate **TODO**
- Be configurable **TODO**

4.1 Bochs

Bochs [13] **TODO**

4.2 Xen

Xen [15] **TODO** *

4.3 Qemu

Qemu [4] **TODO**

4.4 Simics

Wind River Simics [12] **TODO**

5 CPUID Definition and Generation

The described approach
It has the following advantages

- Uses natural unit of configuration state — a field.
- In the meantime, it allows its users to operate in terms of 32 bit leaves values which are more convenient and compact in many cases.
-

6 Evaluation

6.1 Compatibility with existing code

“legacy”

6.2 Extensions

TODO Field flags, such as “hidden”.

7 Conclusions

In this paper we described our approach to simulation of a single but complex processor instruction CPUID .

Model specific registers. Besides CPUID , processors of Intel architecture may have a quite extensive of model specific registers (MSRs), which can be read/written with RDMSR/WRMSR instructions, and also contain bits of processor identification. The simulation of MSRs is also an important aspect, but it is out of this paper focus. Other architectures may have similar facilities, e.g. special purpose registers (SPRs) in PowerPC.

Acknowledgements

Thanks to my parents for raising an awesome me.

References

- [1] Advanced Micro Devices. *AMD64 Architecture Programmer's Manual Volume 1: Application Programming*, 2012.
- [2] ARM Limited. *Application Note 99. Core Type & Revision Identification*, November 2003.
- [3] ARM Limited. *CPUID Base Register*, 2008.
- [4] Fabrice Bellard. QEMU, a Fast and Portable Dynamic Translator. In *FREENIX Track: 2005 USENIX Annual Technical Conference*, 2005.
- [5] Robert R. Collins. CPUID algorithm wars. *Dr. Dobbs*, November 1996.
- [6] Per Fremstad, Wolfgang Fries, Marian Gasparovic, Parwez Hamid, Brian Hatfield, Dick Jorna, Fernando Nogal, and Karl-Erik Stenfors. *IBM System z10 Enterprise Class Technical Guide*. IBM, November 2009.
- [7] IBM. *PowerPC® Microprocessor Family: The Programming Environments Manual for 64-bit Microprocessors Version 3.0*, July 2005.
- [8] Imagination Technologies. *MIPS32 Architecture*, 2014.
- [9] Intel Corporation. *Intel® 64 and IA-32 Architectures Software Developer's Manual. Volumes 1–3*, 2012.
- [10] Intel Corporation. *Intel® Itanium® Architecture Software Developer's Manual. Volumes 1–4*, 2010.
- [11] intruders. ARM CPUID - processor model detection, 2009.
- [12] Peter S. Magnusson, Magnus Christensson, Jesper Eskilson, Daniel Forsgren, Gustav Hallberg, Johan Hogberg, Fredrik Larsson, Andreas Moestedt, and Bengt Werner. Simics: A Full System Simulation Platform. *Computer*, 35:50–58, February 2002.
- [13] Darek Mihoka and Stanislav Shwartsman. Virtualization Without Direct Execution or Jitting: Designing a Portable Virtual Machine Infrastructure. *ISCA-35 Proceedings of the 1st Workshop on Architectural and Microarchitectural Support for Binary Translation*, 2008.
- [14] MIPS Technologies, Inc. *MIPS® Architecture For Programmers Volume I-A: Introduction to the MIPS32® Architecture*, March 2011.
- [15] Ian Pratt. Xen and the art of virtualization, 2006.
- [16] Grigory Rechistov. A set of CPU identification tools for Intel IA-32 and IA-64 systems, 2014.
- [17] D.L. Weaver, T. Germond, and SPARC International. *The SPARC architecture manual: version 9*. PTR Prentice Hall, 1994.