

CS157 Lecture #13

UNION-FIND

UNION OF SETS: n sets, 1 elem each (tree structure) w/ rank 0

- determine which set an element is in by tracing up to its root (find)
- merge two sets using the following rule (union):
 - * if $\text{rank}(u) < \text{rank}(v)$, make u a child of v
 - * if $\text{rank}(u) = \text{rank}(v)$, make 1 a child of the other & increase rank by 1.

seems to support logarithmic time \rightarrow but we can do better!



every time you call find, make each node you encounter a child of the root

\hookrightarrow "path-compression"

ranks no longer correspond with depth of the tree. only follows the rule:

- * if $\text{rank}(u) < \text{rank}(v)$, make u a child of v
- * if $\text{rank}(u) = \text{rank}(v)$, make one a child of the other and increment rank

\rightarrow analyze n operations

HOW DO WE ANALYZE THIS?

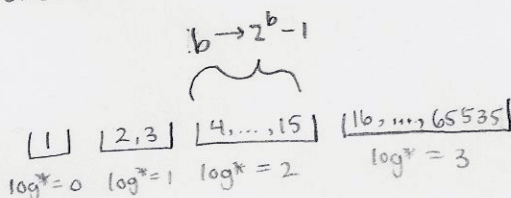
PUNCHLINE: $O(n \log^* n)$, where \log^* is the #times you need to take log to get to 0 or 1 (for all practical purposes, $\log^* n \leq 5$)

* CRUCIAL PROPERTY: every time we process v , rank of $p(v)$ increases [v not root, child of root]

* BUCKETS

* if a node's parent is in the next bucket

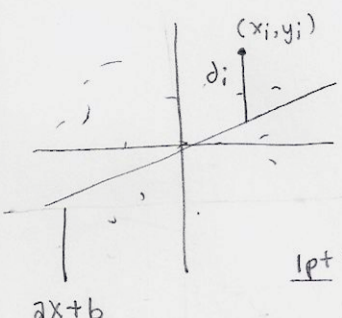
* otherwise, since the rank of its parent always increases, we can do at most $2^{b-1} - b = O(2^b)$ work



$$* \sum_{i=b}^{2^b-1} \frac{n}{2^i} 2^b < \sum_{i=b}^{\infty} \frac{n}{2^i} 2^b = \frac{n}{2^{b+1}} 2^b + \frac{n}{2^{b+2}} 2^b + \dots = n + \frac{n}{2} + \frac{n}{4} + \dots = 2n$$

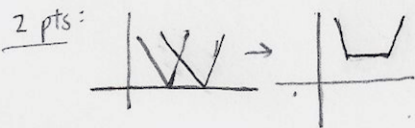
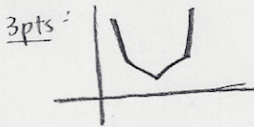
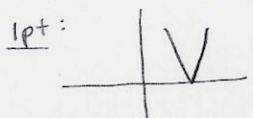
* elem of max rank i parent upgrades

OPTIMIZATION



GOAL: minimize $\sum_i |d_i|$
 $|d_i| = |ax_i + b - y_i|$

* fix a . how do you find b ?



function decreases to a point and then increases
 \hookrightarrow binary search:

* start at a point, and look at the point to the right. if that point has value $<$ curr, recur on it, otherwise, work left

* what about 2 dimensions (a, b unknown)
 \hookrightarrow each col / row looks like previous case
 * choose a pt & one right & up