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#### **Interview Questions: Union-Find**

Help Center

**Warning:** The hard deadline has passed. You can attempt it, but **you will not get credit for it**. You are welcome to try it as a learning exercise.

These interview questions are for your own enrichment and are not assessed. If you click the *Submit Answers* button, you will get a hint.

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# **Question 1**

**Social network connectivity.** Given a social network containing N members and a log file containing M timestamps at which times pairs of members formed friendships, design an algorithm to determine the earliest time at which all members are connected (i.e., every member is a friend of a friend of a friend ... of a friend). Assume that the log file is sorted by timestamp and that friendship is an equivalence relation. The running time of your algorithm should be  $M\log N$  or better and use extra space proportional to N.

### **Question 2**

**Union-find with specific canonical element.** Add a method find() to the union-find data type so that find(i) returns the largest element in the connected component containing i. The operations, union(), connected(), and find() should all take logarithmic time or better.

For example, if one of the connected components is  $\{1,2,6,9\}$ , then the find () method should return 9 for each of the four elements in the connected components because 9 is larger 1, 2, and 6.

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# **Question 3**

Successor with delete. Given a set of N integers  $S=\{0,1,\dots,N-1\}$  and a sequence of requests of the following form:

- ullet Remove x from S
- Find the successor of x: the smallest y in S such that  $y \ge x$ .

design a data type so that all operations (except construction) should take logarithmic time or better.

## **Question 4**

**Union-by-size.** Develop a union-find implementation that uses the same basic strategy as weighted quick-union but keeps track of tree height and always links the shorter tree to the taller one. Prove a  $\lg N$  upper bound on the height of the trees for N sites with your algorithm.

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**Submit Answers** 

Save Answers

You cannot submit your work until you agree to the Honor Code. Thanks!