Data Structure

Tree Part II

Traversal

- Systematic way of visiting/processing all the nodes
- Methods: Preorder, Inorder, and Postorder
- They all traverse the <u>left subtree</u> before the <u>right</u> <u>subtree</u>. It's all about the **position of the root**.



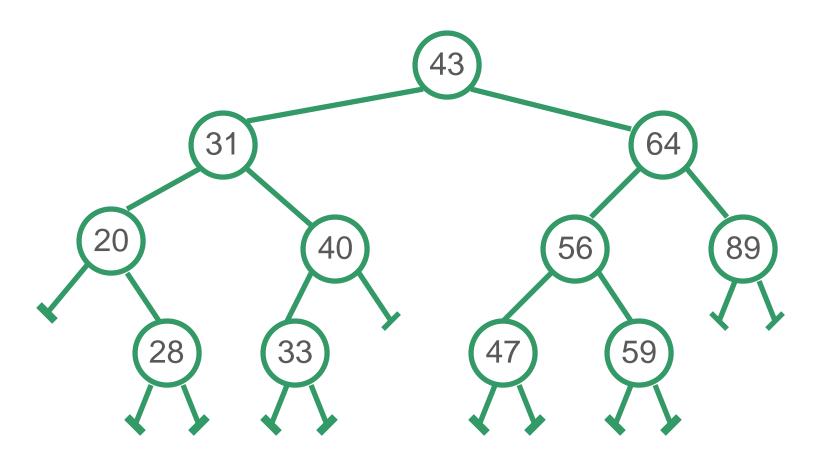




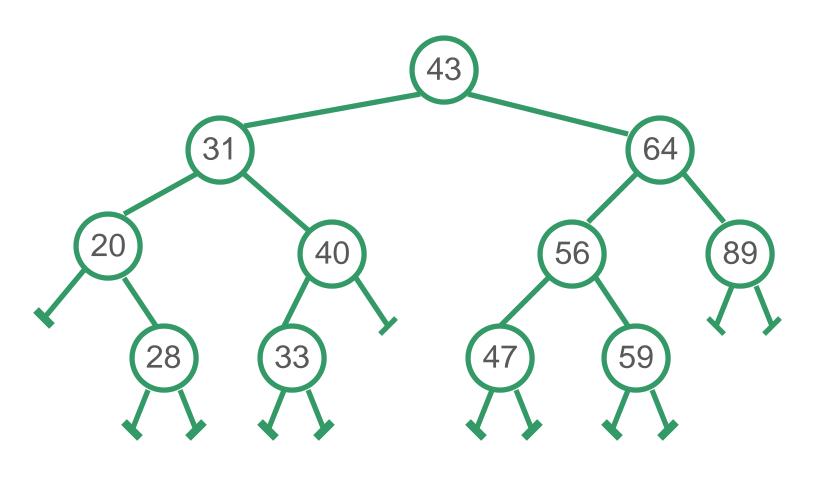




Example: Preorder



Example: Preorder



43 31 20 28 40 33 64 56 47 59 89

- 1) Print the **root** node
- 2) Traverse the **left** subtree
- 3) Traverse the **right** subtree

```
def print_preorder(self):
```

- 1) Print the **root** node
- 2) Traverse the **left** subtree
- 3) Traverse the **right** subtree

```
def print_preorder(self):
    self._print_preorder_aux(self.root)
```

The helper method receives a reference to the "next root"

- 1) Print the **root** node
- 2) Traverse the **left** subtree
- 3) Traverse the **right** subtree

```
def print_preorder(self):
    self._print_preorder_aux(self.root)

def _print_preorder_aux(self, current):
```

- 1) Print the **root** node
- 2) Traverse the **left** subtree
- 3) Traverse the **right** subtree

```
def print_preorder(self):
    self._print_preorder_aux(self.root)

def _print_preorder_aux(self, current):
    if current is not None: # if not a base case
```

Work to do...

- 1) Print the **root** node
- 2) Traverse the **left** subtree
- 3) Traverse the **right** subtree

```
class TreeNode:
    def __init__(self, item=None, left=None, right=None):
        self.item = item
        self.left = left
        self.right = right

def __str__(self):
    return str(self.item)
```

```
def print_preorder(self):
    self._print_preorder_aux(self.root)

def _print_preorder_aux(self, current):
    if current is not None: # if not a base case
        print(current)
```

- 1) Print the **root** node
- 2) Traverse the **left** subtree
- 3) Traverse the **right** subtree

```
def print_preorder(self):
    self._print_preorder_aux(self.root)

def _print_preorder_aux(self, current):
    if current is not None: # if not a base case
        print(current)
        self._print_preorder_aux(current.left)
```

- 1) Print the **root** node
- 2) Traverse the **left** subtree
- 3) Traverse the **right** subtree

```
def print_preorder(self):
    self._print_preorder_aux(self.root)

def _print_preorder_aux(self, current):
    if current is not None: # if not a base case
        print(current)
        self._print_preorder_aux(current.left)
        self._print_preorder_aux(current.right)
```

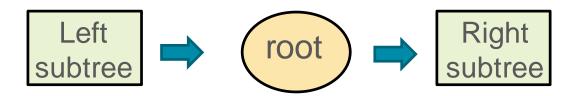
```
def print_preorder(self):
    self._print_preorder_aux(self.root)

def _print_preorder_aux(self, current):
    if current is not None: # if not a base case
        print(current)
        self._print_preorder_aux(current.left)
        self._print_preorder_aux(current.right)
```

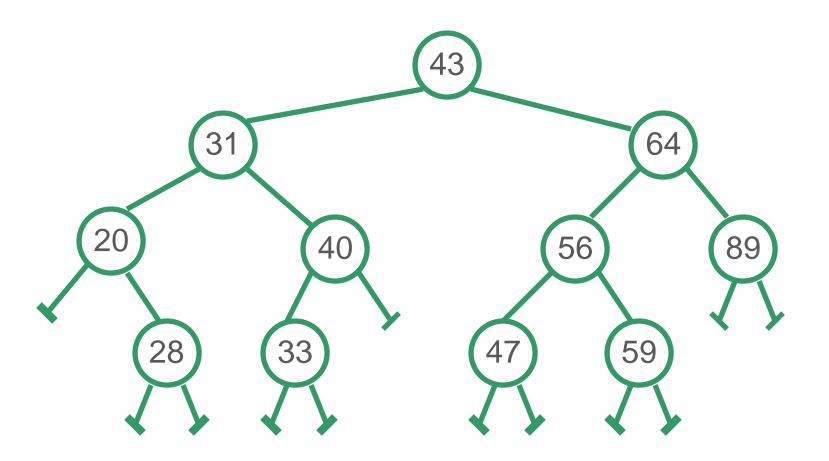


Print Inorder Traversal

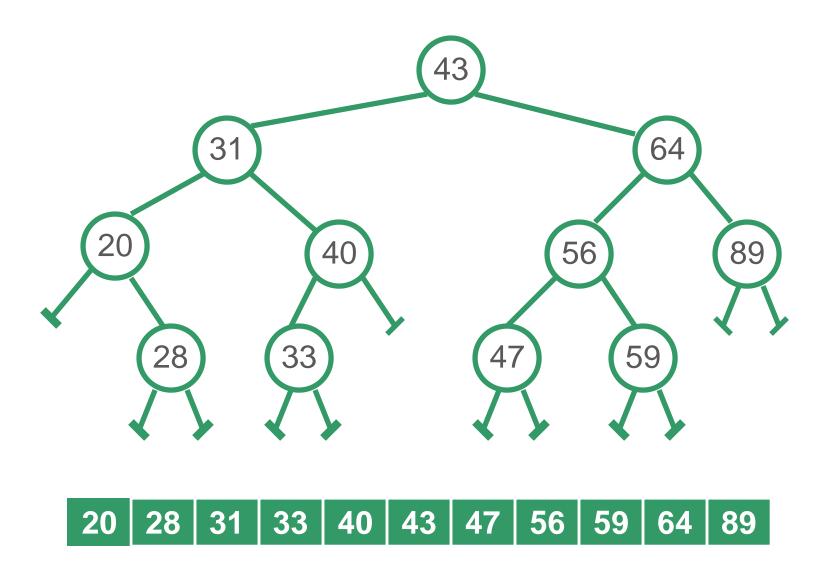
Inorder



Example: Inorder



Example: Inorder





Print In-order Traversal

- 1) Traverse the **left** subtree
- 2) Print the root node
- 3) Traverse the **right** subtree

```
def print_inorder(self):
    self._print_inorder_aux(self.root)

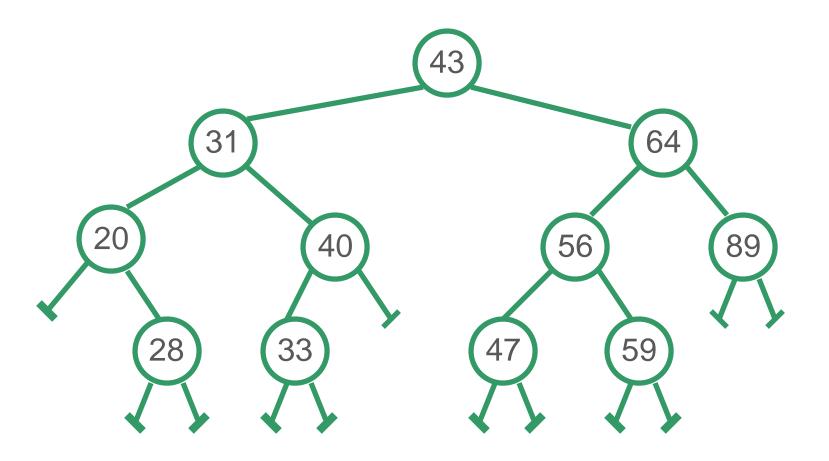
def _print_inorder_aux(self, current):
    if current is not None: # if not a base case
        self._print_inorder_aux(current.left)
        print(current)
        self._print_inorder_aux(current.right)
```

Print Postorder Traversal

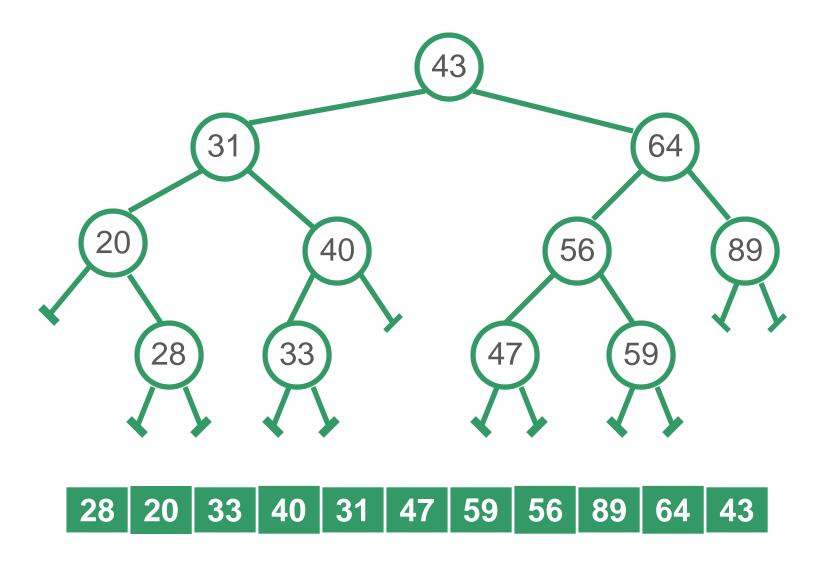
Postorder



Example: Postorder



Example: Postorder





Print Post-order Traversal

- 1) Traverse the **left** subtree
- 2) Traverse the **right** subtree
- 3) Print the **root** node

```
def print_postorder(self):
    self._print_postorder_aux(self.root)

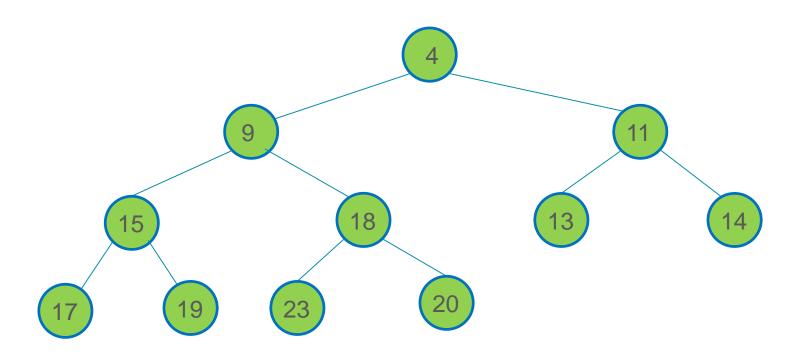
def _print_postorder_aux(self, current):
    if current is not None: # if not a base case
        self._print_postorder_aux(current.left)
        self._print_postorder_aux(current.right)
        print(current)
```

Heap

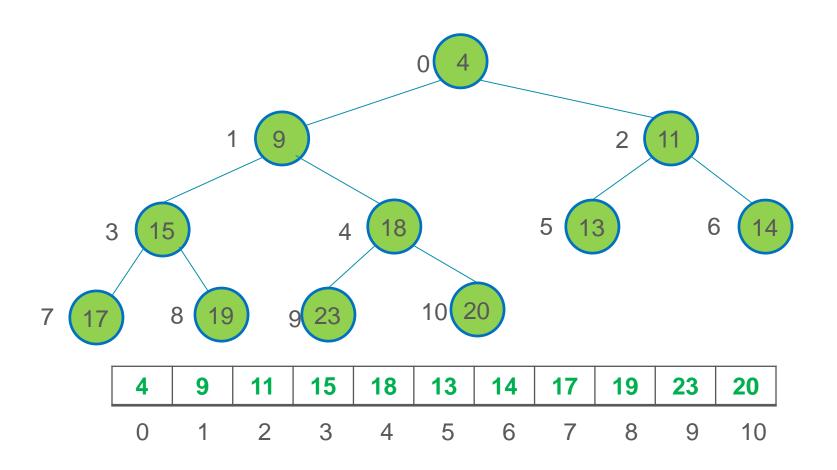
- A specialized tree-based data structure which is a complete binary tree.
- There are two types:
- Max-heap
- Min-heap

Min-Heap

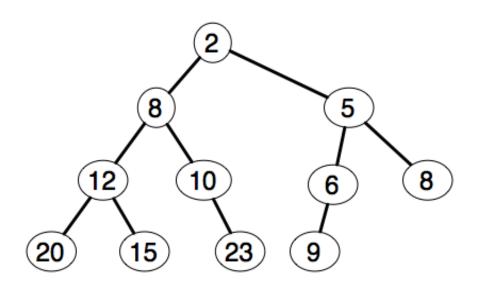
- A min-heap is a balanced binary tree in which:
- ✓ each parent is always smaller than or equal to its children (this implies that the root is the smallest element in the heap)



Heap can be represented as List



The following tree is a min-Heap.



- A. Yes
- B. No

Consider a min-heap represented by the list shown below, which of the following is false?

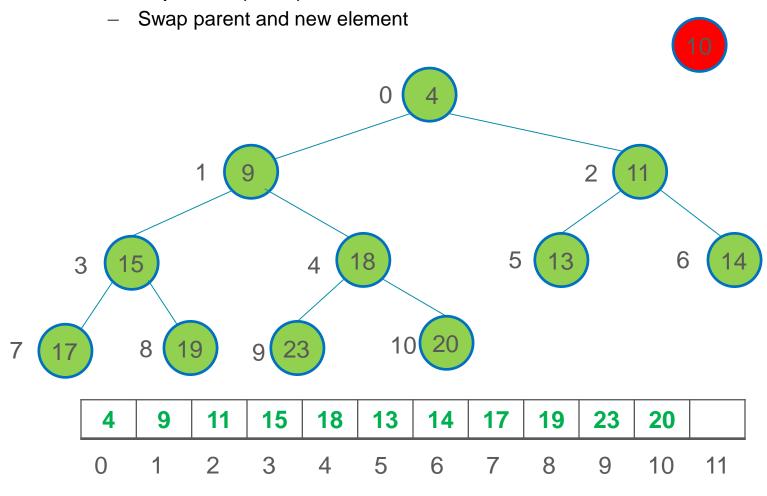
- A. E is a right child of C
- B. D is a left child of C
- C. H is a parent of L
- D. D is a parent of K
- E. A is a parent of C

Α	В	С	F	Н	D	Е	G	J	L	K
0	1	2	3	4	5	6	7	8	9	10

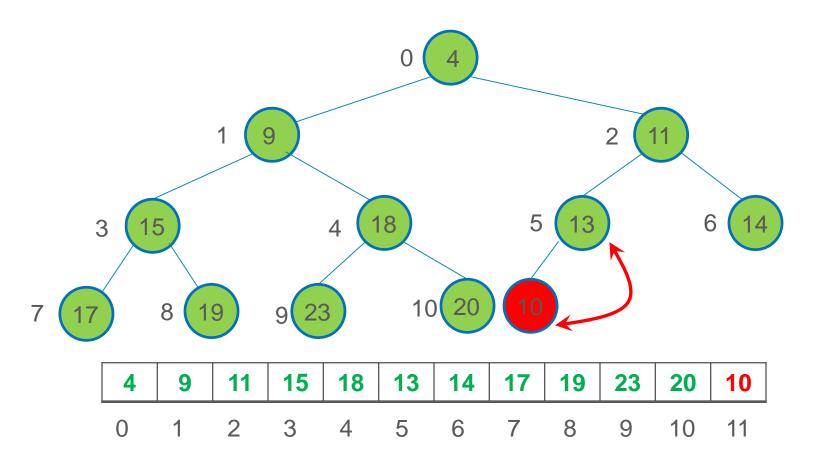
Min-Heap Operations

- Insert: insert a value in the min-heap and ensure that heap properties are maintained
- Extraction: Retrieve and delete the minimum value from heap and ensure that heap properties are maintained

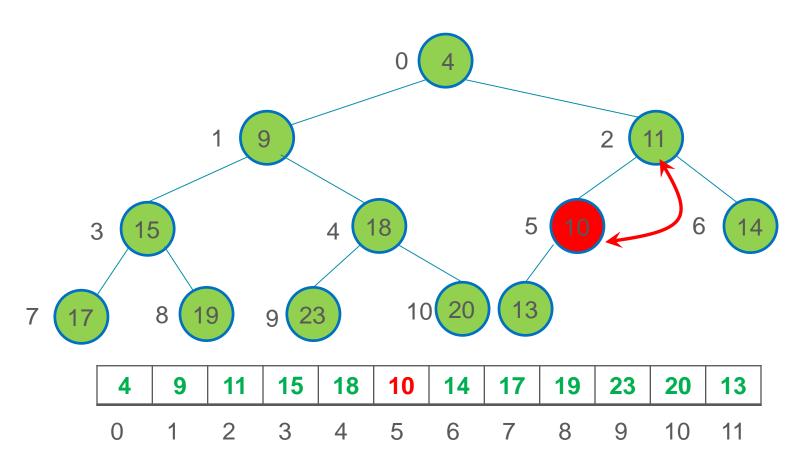
- Append the new element in hList
- While parent(new) > new



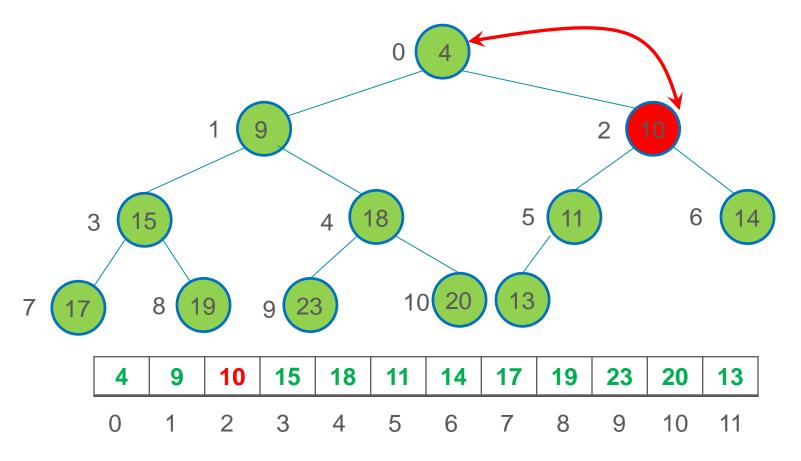
- Append the new element in hList
- While parent(new) > new
 - Swap parent and new element



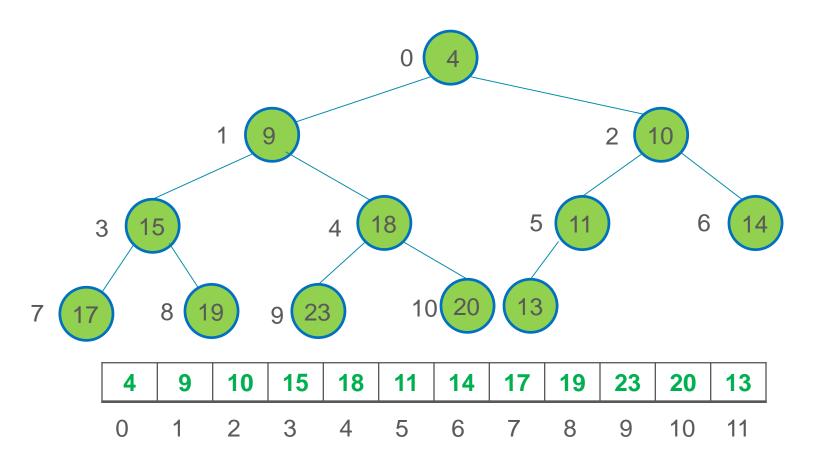
- Append the new element in hList
- While parent(new) > new
 - Swap parent and new element



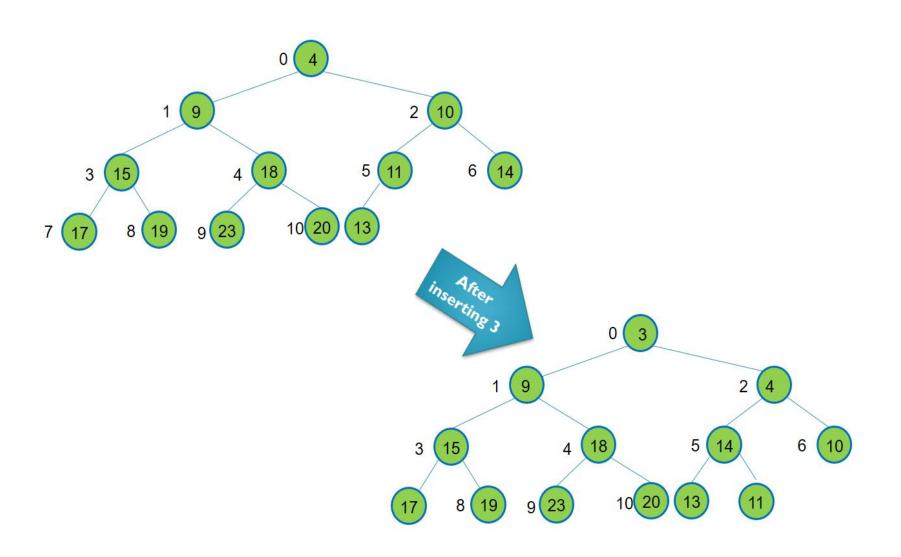
- Append the new element in hList
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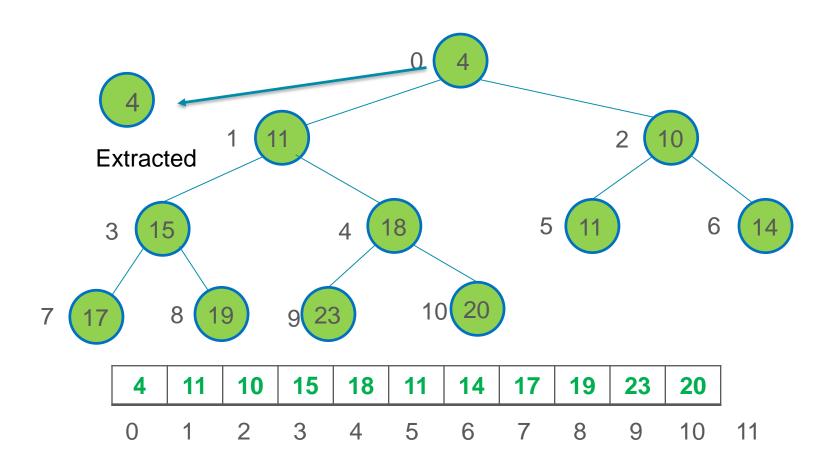
After inserting 3, is the heap updated as shown below?



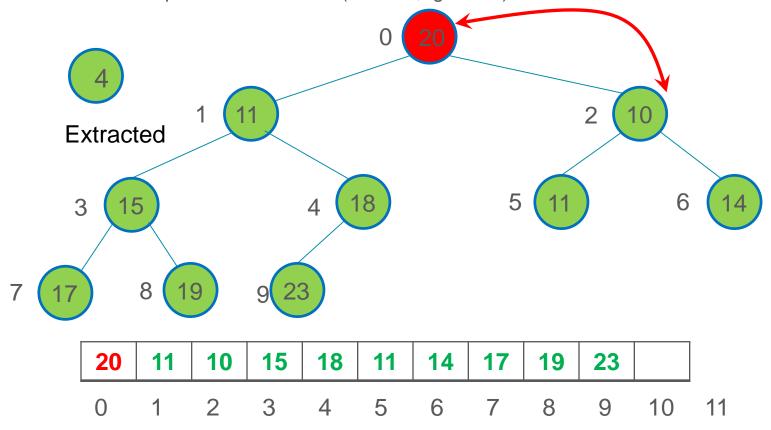
Insertion to Max-heap

- Insert the following items (in the order they are listed) into a maxheap:
- 10, 16, 2, 29, 6, 14, 11, 4.

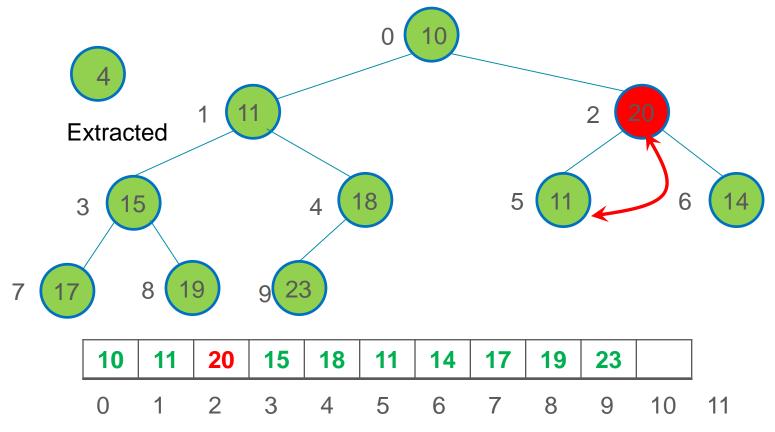
- Store and delete hList[0]
- Move hList[N-1] (called last) to hList[0]



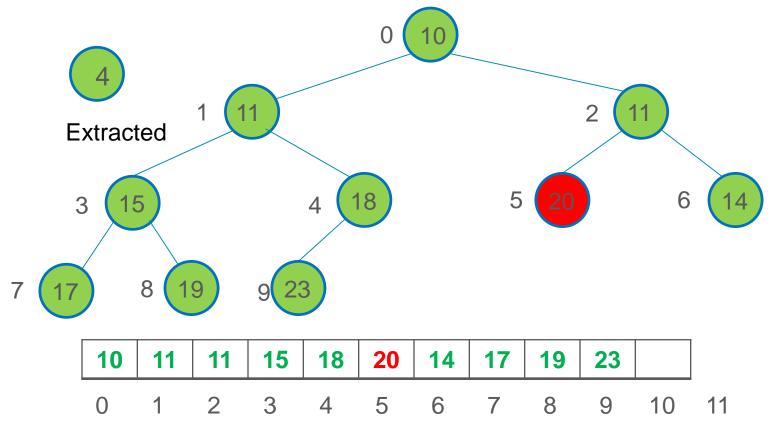
- Store and delete hList[0]
- Move hList[N-1] (called last) to hList[0]
- While leftchild(last) < last or rightchild(last) < last
 - Swap last with minimum(leftchild, rightchild)



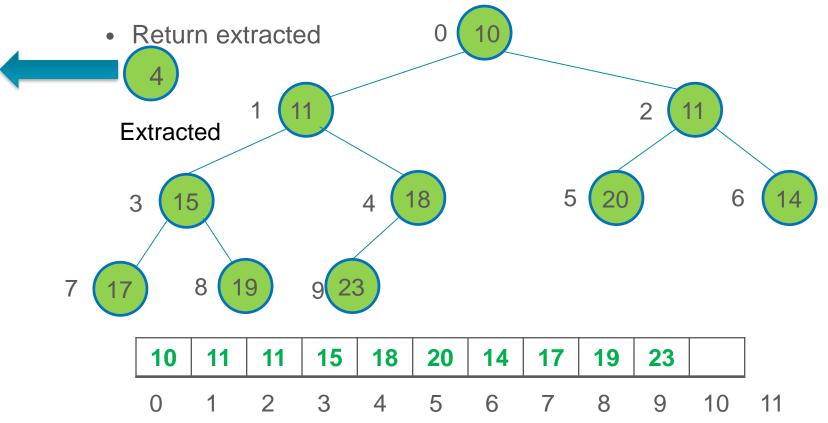
- Store and delete hList[0]
- Move hList[N-1] (called last) to hList[0]
- While leftchild(last) < last or rightchild(last) < last
 - Swap last with minimum(leftchild,rightchild)



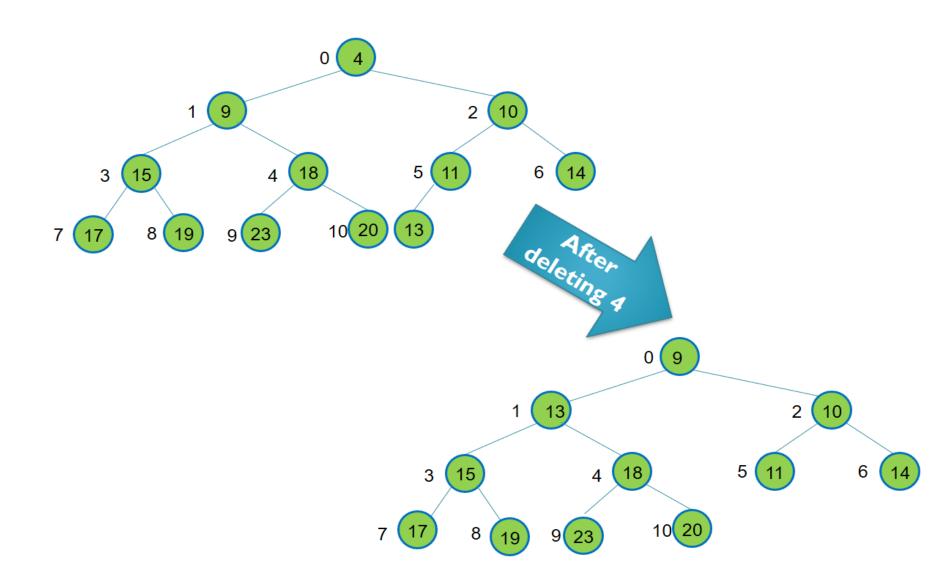
- Store and delete hList[0]
- Move hList[N-1] (called last) to hList[0]
- While leftchild(last) < last or rightchild(last) < last
 - Swap last with minimum(leftchild,rightchild)



- Store and delete hList[0]
- Move hList[N-1] (called last) to hList[0]
- While leftchild(last) < last or rightchild(last) < last
 - Swap last with minimum(leftchild,rightchild)



After deleting 4, is the heap updated as shown below?



How the **Heap Sort** works?

• Discuss how a **heap** can be used to **sort a list**. Using a heap, sort the list [10, 16, 2, 29, 6, 14, 11, 4] in **ascending order**. Show the state of the heap at every step.

Final Announcement

- No More Lab Session this week, instead there will be a revision session on Thursday at 1:00 pm.
- Lab Effort Mark
- Thank you for your attendance and cooperation
- Please complete LUMES module evaluation
- Thank you and Stay Safe ☺