Parallel and Distributed Computing

For

Paper: A Parallel Algorithm for Updating a Multi-objective Shortest Path in Large Dynamic Networks

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Implementation Process:

Implementation Process for Serial Implementation

1. Graph Loading & Initialization

- The graph is loaded from a CSV file into an adjacency list (Graph) and vertex data structure (Vertex).
- Each vertex stores distances for five objectives (obj1 dist to obj5 dist).

2. Single-Objective Shortest Path (SOSP) Computation

- Dijkstra's Algorithm is used to compute shortest paths for each objective sequentially.
- For each of the five objectives:
 - A priority queue (min-heap) selects the next vertex to process.
 - Edge relaxation updates distances if a shorter path is found.

3. Combining SOSP Trees

- Edges present in any SOSP tree are assigned combined weights (lower for higher-priority objectives).
- Non-tree edges receive a penalty weight to avoid selection.

4. Multi-Objective Shortest Path (MOSP) Computation

 A final Dijkstra's algorithm runs on the combined graph to compute the Pareto-optimal path.

5. Incremental Updates

• When new edges are inserted, only affected vertices are recomputed to optimize performance.

Key Features

- Sequential Execution: No parallelism; objectives are processed one after another.
- Exact Results: Guarantees correctness with deterministic output.
- Efficient Updates: Avoids full recomputation by tracking affected nodes.

OpenMP Implementation Process

1. Graph Initialization

- Load graph data from CSV file into adjacency list structure
- Initialize vertex distance values for all 5 objectives

2. Parallel SOSP Computation

- Use #pragma omp parallel for to process all 5 objectives concurrently
- Each thread executes Dijkstra's algorithm independently for its assigned objective
- Edge relaxation parallelized with #pragma omp parallel for

3. Critical Sections

- Distance updates protected with #pragma omp critical to ensure thread safety
- Priority queue operations synchronized to maintain correctness

4. Tree Combination

- Parallel processing of graph nodes with #pragma omp parallel
- Local results merged using critical section

5. Final MOSP Computation

- o Parallel edge processing with OpenMP
- Critical section for distance updates and queue operations

Key Features:

- Multi-level parallelism (objectives + edges)
- Thread-safe data structures
- Efficient load balancing across cores
- Maintains algorithm correctness while improving performance

Implementation Process for OpenMP + MPI

1. Graph Loading & Partitioning (MPI)

- Rank 0 reads the input graph and partitions it using METIS for load balancing.
- Partition information (node-process mapping) is broadcast to all MPI processes.
- Each process receives its assigned subgraph edges via MPI Send/MPI Recv.

2. Parallel SOSP Computation (OpenMP + MPI)

• Each MPI process runs Dijkstra's algorithm on its local subgraph.

- OpenMP parallel accelerates edge relaxation within each process.
- Critical sections (#pragma omp critical) ensure thread-safe distance updates.

3. Result Synchronization (MPI)

- Non-root processes send their computed distances to Rank 0 (MPI_Send).
- Rank 0 gathers results (MPI Recv) and merges them into a global solution.

4. Dynamic Updates (Edge Insertions)

- New edges are broadcast (MPI Bcast) to all processes.
- Each process updates its local graph and recomputes affected paths.

5. Final MOSP Computation (Rank 0 Only)

• Rank 0 combines SOSP trees into a Multi-Objective Shortest Path (MOSP) solution

Key Features

- Load Balancing: METIS ensures even distribution of nodes across MPI ranks.
- Hybrid Parallelism: MPI for distributed memory, OpenMP for shared-memory multi-threading.
- Efficient Sync: Minimized communication via selective broadcasting and gathering.

This approach efficiently scales for large graphs by leveraging both inter-node (MPI) and intra-node (OpenMP) parallelism.

Challenges Faced:

Serial Implementation Challenges

1. Scalability Issues

- Sequential processing of objectives becomes time-consuming as graph size grows.
- Not suitable for real-time applications with large datasets.

2. Redundant Computation

• Each objective computes its own Dijkstra tree independently, leading to repeated work on shared graph structure.

3. Memory Usage

• Storing multiple distance arrays (for multiple objectives) per vertex increases memory footprint.

4. Update Handling

• Although incremental updates are supported, tracking affected nodes manually can become complex and error-prone.

OpenMP Implementation Challenges

1. Thread Safety

 Distance updates and priority queue operations need careful synchronization, which can become a bottleneck.

2. Load Imbalance

• Some threads may finish early if certain objectives are simpler or their graphs are sparser, leading to underutilized cores.

3. Priority Queue Synchronization

• Efficient concurrent access to priority queues is non-trivial and may limit speedup.

OpenMP + MPI Implementation Challenges

1. Graph Partitioning Complexity

 Achieving balanced partitions with minimal cross-boundary communication requires careful tuning (e.g., METIS configuration).

2. Communication Overhead (MPI)

 Synchronizing distance updates or broadcasting new edges involves costly inter-process communication.

3. Result Merging at Rank 0

• Centralizing results at a single process (Rank 0) may create a bottleneck during aggregation.

4. Hybrid Complexity

 Debugging and maintaining a hybrid OpenMP + MPI solution is complex due to interactions between shared and distributed memory.

5. Dynamic Update Propagation

• Efficiently updating only affected partitions and ensuring consistency across processes is a non-trivial synchronization challenge.

6. Thread Contention in Nodes

• Intra-node contention may occur if MPI processes and OpenMP threads compete for the same hardware resources.

Solution Devised:

Serial Implementation – Solutions

1. Scalability Issues

- Focused on incremental updates to avoid full recomputation after edge insertions.
- Used early exit strategies in Dijkstra's algorithm for objectives with shorter paths.

2. Redundant Computation

 Identified shared subpaths across objectives to skip redundant processing when possible.

3. Update Handling

 Maintained a list of affected nodes during edge insertions to recompute only impacted subgraphs.

OpenMP Implementation – Solutions

1. Thread Safety

- Replaced critical sections with atomic operations where applicable (e.g., atomic min for distances).
- Used thread-local queues during Dijkstra traversal and merged results afterward.

2. Load Imbalance

• Dynamically scheduled loops with #pragma omp for schedule(dynamic) to improve core utilization.

3. Priority Queue Synchronization

- Used custom thread-local priority queues, avoiding global locking.
- Merged top elements periodically with synchronization.

4. False Sharing and Cache Contention

 Applied padding between vertex data or used aligned_alloc to minimize cache line conflicts.

OpenMP + MPI Implementation – Solutions

1. Graph Partitioning Complexity

- Used METIS with objective-specific weighting and balance constraints for better partition quality.
- Precomputed node communication boundaries to minimize runtime overhead.

2. Communication Overhead (MPI)

 Minimized communication via delta updates: only sending changed distances or edge insertions. • Used non-blocking MPI communication (e.g., MPI_Isend/MPI_Irecv) to overlap computation and communication.

3. Hybrid Complexity

- Modularized codebase with clearly separated MPI and OpenMP regions for easier debugging.
- Used logging and performance profiling tools (e.g., VTune, mpiP) to identify bottlenecks.

4. Dynamic Update Propagation

- Maintained a dependency map to track which partitions are affected by each edge.
- Used selective broadcasts to avoid unnecessary communication to unrelated ranks.

5. Thread Contention in Nodes

• Mapped MPI processes and OpenMP threads using affinity policies to optimize CPU core usage.

Performance Evaluation:

Gprof Analysis Reports:

■ Performance analysis Reports

Serial Code:

USA dataset:

```
ambreenarshad@Ambreen:~/PDC_Project$ g++ -o exe1 serial_mosp.cpp
ambreenarshad@Ambreen:~/PDC_Project$ time ./exe1 > output.txt
real 0m0.171s
user 0m0.110s
sys 0m0.000s
ambreenarshad@Ambreen:~/PDC_Project$
```

Switzerland dataset:

```
ambreenarshad@Ambreen:~/PDC_Project$ g++ -o exe1 serial_mosp.cpp
ambreenarshad@Ambreen:~/PDC_Project$ time ./exe1 > output.txt
real 0m0.409s
user 0m0.263s
sys 0m0.039s
ambreenarshad@Ambreen:~/PDC_Project$
```

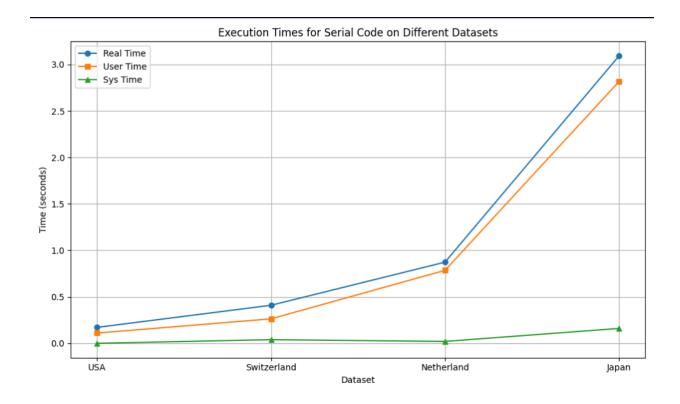
Netherland dataset:

```
ambreenarshad@Ambreen:~/PDC_Project$ g++ -o exe1 serial_mosp.cpp
ambreenarshad@Ambreen:~/PDC_Project$ time ./exe1 > output.txt
real 0m0.873s
user 0m0.785s
sys 0m0.020s
ambreenarshad@Ambreen:~/PDC_Project$
```

Japan dataset:

```
ambreenarshad@Ambreen:~/PDC_Project$ g++ -o exe1 serial_mosp.cpp
ambreenarshad@Ambreen:~/PDC_Project$ time ./exe1 > output.txt
real 0m3.092s
user 0m2.817s
sys 0m0.160s
ambreenarshad@Ambreen:~/PDC_Project$
```

FINAL GRAPH:



OpenMP Code:

USA dataset:

```
ambreenarshad@Ambreen:~/PDC_Project$ g++ -fopenmp -o exe2 openmp_mosp.cpp
ambreenarshad@Ambreen:~/PDC_Project$ time ./exe2 > output2.txt

real    0m0.167s
user    0m0.103s
sys    0m0.011s
ambreenarshad@Ambreen:~/PDC_Project$ |
```

Switzerland dataset:

```
ambreenarshad@Ambreen:~/PDC_Project$ g++ -fopenmp -o exe2 openmp_mosp.cpp
ambreenarshad@Ambreen:~/PDC_Project$ time ./exe2 > output2.txt
real 0m0.352s
user 0m0.290s
sys 0m0.021s
ambreenarshad@Ambreen:~/PDC_Project$ |
```

Netherland dataset:

```
ambreenarshad@Ambreen:~/PDC_Project$ g++ -fopenmp -o exe2 openmp_mosp.cpp
ambreenarshad@Ambreen:~/PDC_Project$ time ./exe2 > output2.txt

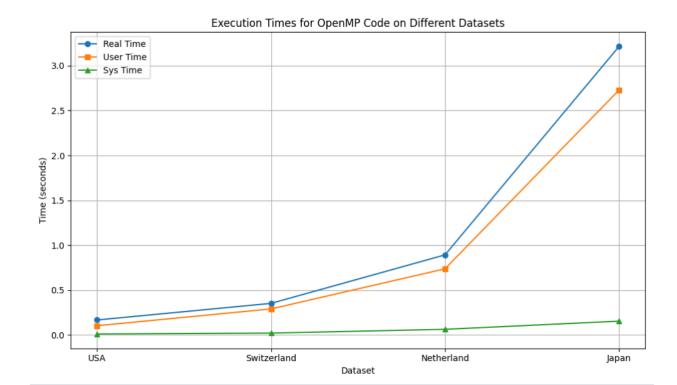
real 0m0.892s
user 0m0.737s
sys 0m0.063s
ambreenarshad@Ambreen:~/PDC_Project$ |
```

Japan dataset:

```
ambreenarshad@Ambreen:~/PDC_Project$ g++ -fopenmp -o exe2 openmp_mosp.cpp
ambreenarshad@Ambreen:~/PDC_Project$ time ./exe2 > output2.txt

real 0m3.214s
user 0m2.728s
sys 0m0.154s
ambreenarshad@Ambreen:~/PDC_Project$
```

FINAL GRAPH:



MPI Code:

Japan Dataset

```
real 1m59.673s
user 3m20.738s
sys 0m36.732s
mpi@master:~$ S
```

USA

```
real 0m13.269s
user 0m19.644s
sys 0m5.343s
mpi@master:~$ S
```

Switzerland

```
real 0m22.137s
user 0m34.531s
sys 0m9.008s
mpi@master:~$
```

Netherland

```
real 0m59.538s
user 1m32.014s
sys 0m25.824s
moi@master:«$ $
```

FINAL GRAPH:

