

Representations of Lattices

Athar Abdul-Quader

Purchase College, SUNY

CUNY Models of Peano Arithmetic Seminar
April 2, 2024

Substructure / Interstructure Lattices

- 1 If $\mathcal{M} \models \text{PA}$, then $\text{Lt}(\mathcal{M}) = (\{\mathcal{K} : \mathcal{K} \prec \mathcal{M}\}, \preceq)$ is a lattice, called the **substructure lattice** of \mathcal{M} .
- 2 If $\mathcal{M} \preceq \mathcal{N}$, then $\text{Lt}(\mathcal{N}/\mathcal{M}) = (\{\mathcal{K} : \mathcal{M} \preceq \mathcal{K} \preceq \mathcal{N}\}, \preceq)$ is a lattice, called the **interstructure lattice** between \mathcal{M} and \mathcal{N} .

Question

For which finite lattices L is there \mathcal{M} such that $\text{Lt}(\mathcal{M}) \cong L$?

Similarly for interstructure lattices: given (countable / nonstandard) \mathcal{M} , for which finite lattices L is there $\mathcal{M} \prec \mathcal{N}$ such that $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong L$?

Schmerl

Many results are known (e.g. \aleph_1 -algebraic distributive lattices), but the question remains open.

We are also interested in the nature of these extensions. That is: given a (countable, nonstandard) \mathcal{M} , and a finite L , is there $\mathcal{M} \prec_{\text{end}} \mathcal{N}$ such that $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong L$? Can such an extension be conservative? Cofinal? Mixed?

Many results rely on a technique devised by Schmerl (1986): **representations of lattices**. From now on, we will restrict our attention to countable models.

Outline

For today, we will study the Boolean Algebra \mathbf{B}_2 to motivate this technique.

- 1 First, we will consider $\mathcal{M} \prec \mathcal{M}(c)$ where $\text{Lt}(\mathcal{M}(c)/\mathcal{M}) \cong \mathbf{B}_2$. We will recover certain information about $\text{tp}(c/\mathcal{M})$ based on the fact that the interstructure lattice is \mathbf{B}_2 .
- 2 This will motivate the definition of representations of lattices, along with properties of those representations.
- 3 We will take a countable \mathcal{M} and show it has a set of representations of \mathbf{B}_2 satisfying some of these properties, using this to construct a type $p(x)$ which, when realized, will generate a model \mathcal{N} where $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong \mathbf{B}_2$.
- 4 This, in turn, motivates the main result, due to Schmerl, showing the connection between realizing a lattice as an interstructure lattice and the existence of certain representations of that lattice.

The Lattice \mathbf{B}_2

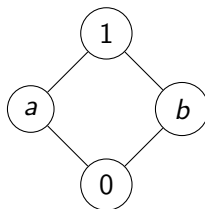
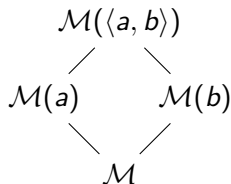


Figure: The lattice \mathbf{B}_2

For now: fix $\mathcal{M} \prec \mathcal{N}$ is such that $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong \mathbf{B}_2$.

Generators

B_2 is finite: so each element is compact. There are $a, b \in N$ such that $\text{Lt}(\mathcal{N}/\mathcal{M})$ is exactly:



In fact: $\text{tp}(\langle a, b \rangle / \mathcal{M})$ **knows** this! (How can we make this precise?)

$tp(\langle a, b \rangle)$

Suppose $c \in N$. Let t be such that $\mathcal{N} \models t(\langle a, b \rangle) = c$. One of the following must hold:

- ① $c \in M$.
- ② $\mathcal{M}(c) = \mathcal{M}(a)$.
- ③ $\mathcal{M}(c) = \mathcal{M}(b)$.
- ④ $\mathcal{M}(c) = \mathcal{N}$.

If $c \in M$: Let

$$X = \{x : t(x) = c\}.$$

X is an infinite, \mathcal{M} -definable set, $\langle a, b \rangle \in X^{\mathcal{N}}$ and $\mathcal{M} \models t \upharpoonright X$ is constant.

$\mathcal{M}(a)$

If $\mathcal{M}(c) = \mathcal{M}(a)$: There are \mathcal{M} -definable g_1, g_2 such that $\mathcal{N} \models g_1(c) = a \wedge g_2(a) = c$. Let π_1 be the projection function $\pi_1(\langle x, y \rangle) = x$. Let

$$X = \{x : g_1(t(x)) = \pi_1(x) \wedge g_2(\pi_1(x)) = t(x)\}.$$

X is an infinite, \mathcal{M} -definable set, $\langle a, b \rangle \in X^{\mathcal{N}}$, and $\mathcal{M} \models \forall x, y \in X (t(x) = t(y) \leftrightarrow \pi_1(x) = \pi_1(y))$.

$\mathcal{M}(b)$

If $\mathcal{M}(c) = \mathcal{M}(b)$: There are \mathcal{M} -definable g_1, g_2 such that $\mathcal{N} \models g_1(c) = b \wedge g_2(b) = c$. Let π_2 be the projection function $\pi_2(\langle x, y \rangle) = y$. Let

$$X = \{x : g_1(t(x)) = \pi_2(x) \wedge g_2(\pi_2(x)) = t(x)\}.$$

X is an infinite, \mathcal{M} -definable set, $\langle a, b \rangle \in X^{\mathcal{N}}$, and $\mathcal{M} \models \forall x, y \in X (t(x) = t(y) \leftrightarrow \pi_2(x) = \pi_2(y))$.

\mathcal{N}

If $\mathcal{M}(c) = \mathcal{N}$: there is \mathcal{M} -definable g_1 such that $\mathcal{N} \models g_1(c) = \langle a, b \rangle$. Let

$$X = \{x : t(x) = c \wedge g_1(c) = x\}.$$

X is an infinite, \mathcal{M} -definable set, $\langle a, b \rangle \in X^{\mathcal{N}}$, and $\mathcal{M} \models \forall x, y \in X (t(x) = t(y) \leftrightarrow x = y)$.

Notice

Whenever t is an \mathcal{M} -definable (total) function, consider the equivalence relation induced by t : $x \sim y$ iff $t(x) = t(y)$. We can find an infinite, \mathcal{M} -definable X such that $\langle a, b \rangle \in X^{\mathcal{N}}$ and one of the following holds on X :

- ① \sim is trivial (ie, for all $x, y \in X$, $x \sim y$).
- ② $x \sim y$ iff $\pi_1(x) = \pi_1(y)$.
- ③ $x \sim y$ iff $\pi_2(x) = \pi_2(y)$.
- ④ \sim is discrete (ie, for all $x, y \in X$, $x \sim y$ iff $x = y$).

Definitions

Let X be a set and L a finite lattice. Then:

- 1 $\text{Eq}(X)$ is the lattice of equivalence relations on X , with top element $\mathbf{1}_X = X \times X$ (trivial relation) and bottom element $\mathbf{0}_X = \{(a, a) : a \in X\}$ (discrete relation).
- 2 $\alpha : L \rightarrow \text{Eq}(X)$ is a **representation** if it is one to one and:
 - $\alpha(0_L) = \mathbf{1}_X$ ($\alpha(0)$ is trivial),
 - $\alpha(1_L) = \mathbf{0}_X$ ($\alpha(1)$ is discrete), and
 - $\alpha(x \vee y) = \alpha(x) \wedge \alpha(y)$.

That is, a representation picks out specific equivalence relations on X , one for each $r \in L$. Ensure that if $x \leq y$, then $\alpha(y)$ refines $\alpha(x)$.

Example

Let X be the set of pairs $\{\langle x, y \rangle : x < y\}$. Define $\alpha : \mathbf{B}_2 \rightarrow \text{Eq}(X)$:

- $\alpha(0)$ is trivial,
- $(\langle a_1, a_2 \rangle, \langle b_1, b_2 \rangle) \in \alpha(a)$ iff $a_1 = b_1$,
- $(\langle a_1, a_2 \rangle, \langle b_1, b_2 \rangle) \in \alpha(b)$ iff $a_2 = b_2$, and
- $\alpha(1)$ is discrete.

(Clearly a representation.)

More definitions

Let $\alpha : L \rightarrow \text{Eq}(X)$ be a representation. Then:

- 1 If $Y \subseteq X$, then $\alpha|Y : L \rightarrow \text{Eq}(Y)$ is defined by $(\alpha|Y)(r) = \alpha(r) \cap Y^2$ for each $r \in L$.
- 2 If $\Theta \in \text{Eq}(X)$, Θ is **canonical** for α if there is $r \in L$ such that for all $x, y \in X$, $(x, y) \in \Theta$ iff $(x, y) \in \alpha(r)$.

Notice what happened in our analysis of $\text{tp}(\langle a, b \rangle / \mathcal{M})$: for every \mathcal{M} -definable equivalence relation Θ (on some \mathcal{M} -definable X where $\langle a, b \rangle \in X$), there is \mathcal{M} -definable $Y \subseteq X$ such that $\langle a, b \rangle \in Y^{\mathcal{N}}$ and $\Theta \cap Y^2$ is canonical for $\alpha|Y$.

Plan

Let $\mathcal{M} \models \text{PA}$ be countable. Fix X and $\alpha : \mathbf{B}_2 \rightarrow \text{Eq}(X)$ as before. We wish to construct a complete type $p(x)$ such that if c realizes $p(x)$, then $\text{Lt}(\mathcal{M}(c)/\mathcal{M}) \cong \mathbf{B}_2$.

We do this by constructing a sequence $X_0 \supseteq X_1 \supseteq \dots$ of definable sets that will generate $p(x)$; that is,

$$p(x) = \{\phi(x) : \exists n \in \omega (X_n \subseteq \phi(\mathcal{M}))\}.$$

We ensure that $p(x)$ knows that $\text{Lt}(\mathcal{M}(c)/\mathcal{M}) \cong \mathbf{B}_2$: that is, for each $t(x)$, we ensure that there is X_n where (the equivalence relation induced by) t is canonical for $\alpha|_{X_n}$.

Canonical Ramsey's Theorem

Definition

Suppose $\alpha : L \rightarrow \text{Eq}(A)$ and $\beta : L \rightarrow \text{Eq}(B)$. $\alpha \cong \beta$ if there is a bijection $f : A \rightarrow B$ such that whenever $x, y \in A$ and $r \in L$, then $(x, y) \in \alpha(r)$ iff $(f(x), f(y)) \in \beta(r)$.

Start with $X_0 = X = \{\langle x, y \rangle : x < y\}$. Then: for any \mathcal{M} -definable $\Theta \in \text{Eq}(X)$, there is $Y \subseteq X$ such that $\Theta \cap Y^2$ is canonical for $\alpha|Y$ and $\alpha|Y \cong \alpha$. Why?

Theorem (Canonical Ramsey's Theorem (Erdős-Rado 1950))

For any $k \in \mathbb{N}$ and coloring $c : [\mathbb{N}]^k \rightarrow \mathbb{N}$, there is $l \subseteq k$ and an infinite $H \subseteq \mathbb{N}$ such that whenever $\bar{x}, \bar{y} \in [H]^k$, $c(\bar{x}) = c(\bar{y})$ iff $x_i = y_i$ for all $i \in l$.

Generating a type

Enumerate definable equivalence relations as $\Theta_0, \Theta_1, \dots$. Given X_i , use CRT to find $X_{i+1} \subseteq X_i$ for which Θ_i is canonical. Obtain

$X_0 \supseteq X_1 \dots$ which generates $p(x)$

Claim: $p(x)$ is complete.

Suppose $\phi(x)$ is a formula and let Θ be the equivalence relation given by $(x, y) \in \Theta$ iff $\mathcal{M} \models \phi(x) \leftrightarrow \phi(y)$. Notice that Θ has exactly two classes, but $\alpha(a), \alpha(b), \alpha(1)$ all have infinitely many.

Therefore if Θ is canonical on $\alpha|X_i$, it must be trivial (ie $\Theta \cap X_i^2 = \alpha(0) \cap X_i^2$). That is, one of the following must hold:

- 1 $\mathcal{M} \models \forall x(x \in X_i \rightarrow \phi(x))$, or,
- 2 $\mathcal{M} \models \forall x(x \in X_i \rightarrow \neg \phi(x))$.

Ensuring B_2

Let c realize $p(x)$. How do we show that $\text{Lt}(\mathcal{M}(c)/\mathcal{M}) \cong B_2$? Let $c = \langle a, b \rangle$, then we claim this lattice is exactly $\{\mathcal{M}, \mathcal{M}(a), \mathcal{M}(b), \mathcal{M}(c)\}$.

To see this: let $d \in \mathcal{M}(c)$ and t such that $\mathcal{M}(c) \models t(c) = d$. t is made canonical on some X_i .

That is: there is $r \in B_2$ such that for all $x, y \in X_i$, $t(x) = t(y)$ iff $(x, y) \in \alpha(r)$.

Exercise: if $r = 0$, then $d \in \mathcal{M}$, if $r = a$ then $\mathcal{M}(d) = \mathcal{M}(a)$, if $r = b$, then ...

(Exercise: check that $\mathcal{M}(a) \cap \mathcal{M}(b) = \mathcal{M}$.)

Definitions

Definition (Schmerl 1986)

Let $\alpha : L \rightarrow \text{Eq}(X)$ be a representation.

- 1 α has the *0-canonical partition property*, or is *0-CPP*, if for each $r \in L$, $\alpha(r)$ does not have exactly two classes.
- 2 α is $(n+1)$ -*CPP* if, for each $\Theta \in \text{Eq}(X)$ there is $Y \subseteq X$ such that $\alpha|_Y$ is an n -CPP representation and $\Theta \cap Y^2$ is canonical for $\alpha|_Y$.

The α we used was n -CPP for every $n \in \omega$.

Other properties

We used the following facts about the representations:

- 1 Each $\alpha|X_i$ is 0-CPP (implies $p(x)$ is complete!)
- 2 Whenever $\Theta \in \text{Eq}(X)$ is \mathcal{M} -definable, there is X_i such that Θ is canonical for $\alpha|X_i$.

Definition

Let $\mathcal{M} \models \text{PA}$ and L a finite lattice.

- 1 $\alpha : L \longrightarrow \text{Eq}(X)$ is an \mathcal{M} -representation if α and X are \mathcal{M} -definable.
- 2 If $X \in \text{Def}(\mathcal{M})$, then by $\text{Eq}^{\mathcal{M}}(X)$ we mean the lattice of \mathcal{M} -definable equivalence relations on X .
- 3 \mathcal{C} is an \mathcal{M} -correct set of representations of L if each \mathcal{C} is a nonempty set of 0-CPP \mathcal{M} -representations of L and whenever $\alpha : L \longrightarrow \text{Eq}(X) \in \mathcal{C}$ and $\Theta \in \text{Eq}^{\mathcal{M}}(X)$, there is $Y \subseteq X$ such that $\alpha|Y \in \mathcal{C}$ and $\Theta \cap Y^2$ is canonical for $\alpha|Y$.

Theorem

Theorem (Schmerl 2024)

Let $\mathcal{M} \models \text{PA}$ and L be a finite lattice. Then:

- 1 If there is \mathcal{N} such that $\mathcal{M} \prec \mathcal{N}$ and $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong L$, then there is an \mathcal{M} -correct set of representations of L .
- 2 If \mathcal{M} is countable and there is an \mathcal{M} -correct set of representations of L , then there is $\mathcal{N} \succ \mathcal{M}$ such that $\text{Lt}(\mathcal{N}/\mathcal{M}) \cong L$.

Ranked lattices

A *ranked lattice* (L, ρ) is a lattice L equipped with a function $\rho : L \rightarrow L$ such that for all x and y in L

- ① $x \leq \rho(x)$;
- ② $\rho(\rho(x)) = \rho(x)$;
- ③ $\rho(x) \leq \rho(y)$ or $\rho(y) \leq \rho(x)$;
- ④ $\rho(x \vee y) = \rho(x) \vee \rho(y)$.

The *rankset* of a ranked lattice (L, ρ) is $\{\rho(x) : x \in L\}$. If $\mathcal{M} \prec \mathcal{N}$, the rank function we use is $\text{rk}(\mathcal{K}) = \overline{\mathcal{K}}$, where $\mathcal{K} \preceq_{\text{cof}} \overline{\mathcal{K}} \preceq_{\text{end}} \mathcal{N}$. Then $\text{Ltr}(\mathcal{N}/\mathcal{M}) = (\text{Lt}(\mathcal{N}/\mathcal{M}), \text{rk})$.

Ranking B_2

What is the rankset for the B_2 extension we constructed?

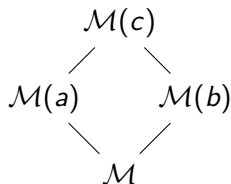


Figure: $\text{Lt}(\mathcal{M}(c)/\mathcal{M}) \cong B_2$

First: $\mathcal{M} \prec_{\text{end}} \mathcal{M}(c)$. (Exercise.)

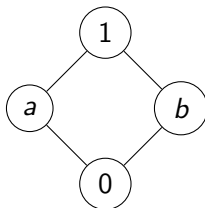
$\mathcal{M}(a)$ and $\mathcal{M}(b)$

Notice: exactly one of $\mathcal{M}(a)$ and $\mathcal{M}(b)$ must be cofinal in $\mathcal{M}(c)$.
Why?

Since $X = \{\langle x, y \rangle : x < y\}$, then $\pi_1(x) < \pi_2(x) \in \text{tp}(c)$, so
 $\mathcal{M}(c) \models a < b$. Hence: $\mathcal{M}(a) \prec_{\text{end}} \mathcal{M}(c)$ and $\mathcal{M}(b) \prec_{\text{cof}} \mathcal{M}(c)$.

The rankset must be $\{\mathcal{M}, \mathcal{M}(a), \mathcal{M}(c)\}$. How did the
representation α imply this?

Other rankings?



What other ranksets could work? How could we realize them?

Definition

Definition (Schmerl 2024)

Let $\mathcal{M} \models \text{PA}$ and (L, ρ) a finite ranked lattice.

- 1 If $A \in \text{Def}(\mathcal{M})$ and $\Theta \in \text{Eq}^{\mathcal{M}}(A)$, a set \mathcal{E} of Θ classes is \mathcal{M} -bounded if there is a bounded $I \in \text{Def}(\mathcal{M})$ such that $I \cap X \neq \emptyset$ for each $X \in \mathcal{E}$.
- 2 $\alpha : L \rightarrow \text{Eq}(A)$ is an \mathcal{M} -representation of (L, ρ) if α is an \mathcal{M} -representation of L and whenever $r \leq s \in L$, $s \leq \rho(r)$ if and only if every $\alpha(r)$ -class is the union of an \mathcal{M} -bounded set of $\alpha(s)$ -classes.
- 3 \mathcal{C} is an \mathcal{M} -correct set of representations of (L, ρ) if \mathcal{C} is an \mathcal{M} -correct set of representations of L and each $\alpha \in \mathcal{C}$ is an \mathcal{M} -correct representation of (L, ρ) .

Theorem

Theorem

Suppose $\mathcal{M} \models \text{PA}$ and (L, ρ) is a finite ranked lattice.

- ① If there is \mathcal{N} such that $\mathcal{M} \prec \mathcal{N}$ and $\text{Ltr}(\mathcal{N}/\mathcal{M}) \cong (L, \rho)$, then there is an \mathcal{M} -correct set of representations of (L, ρ) .*
- ② If \mathcal{M} is countable and there is an \mathcal{M} -correct set of representations of (L, ρ) , then there is $\mathcal{N} \succ \mathcal{M}$ such that $\text{Ltr}(\mathcal{N}/\mathcal{M}) \cong (L, \rho)$.*