|   | Assignment 4  |
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|   | Problem Statement: Write a program using TCP socket for wired network for the following     |
|   | a. Say Hello to Each other  |
|   | b. File transfer  |
|   | c. Calculator (Arithmetic)  |
|   | d. Calculator (Trigonometry)  |
|   | Demonstrate the packets captured traces using Wireshark Packet Analyzer Tool for peer to    |
|   | peer mode.  |
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|   | Objective:  |
| _ | Setup connection TCP between two nodes  |
|   | a. Say Hello to Each other  |
|   | b. File transfer  |
| _ | c. Calculator (Arithmetic)  |
|   | d. Calculator (Trigonometry)  |
|   | Learning Outcome: Students will be able to  |
|   | · Demonstrate working of TCP sockets  |
|   | · Perform the above operations over TCP sockets   |
|   | Requirements:   |
|   | · Open source linux based OS  |
|   | · Eclipse IDE or Python interpreter   |
|   | Theory  |
|   | TCP   |
|   | TCP is connection-oriented, and a connection between client and server is established befor |
|   | data can be sent. The server must be listening (passive open) for connection requests from  |
|   | clients before a connection is established. Three-way handshake (active open), re-          |
|   | transmission, and error-detection adds to reliability but lengthens latency. TCP employs    |

| network congestion avoidance. However, there are vulnerabilities to TCP including denial of          |
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| service, connection hijacking, TCP veto, and reset attack. For network security, monitoring,         |
| and debugging, TCP traffic can be intercepted and logged with a packet sniffer.                      |
| TCP Segment structure  |
| A TCP segment consists of a segment header and a data section. The segment header contains           |
| 10 mandatory fields, and an optional extension field (Options, pink background in table). The        |
| data section follows the header and is the payload data carried for the application. The length      |
| of the data section is not specified in the segment header; It can be calculated by subtracting      |
| the combined length of the segment header and IP header from the total IP datagram length            |
| specified in the IP header.  |
| Source port (16 bits): Identifies the sending port.  |
| Destination port (16 bits): Identifies the receiving port.   |
| Sequence number (32 bits)  |
| Has a dual role:   |
| $\cdot$ If the SYN flag is set (1), then this is the initial sequence number. The sequence number of |
| the actual first data byte and the acknowledged number in the corresponding ACK are then             |
| this sequence number plus 1.   |
| · If the SYN flag is clear (0), then this is the accumulated sequence number of the first data       |
| byte of this segment for the current session.  |
| Acknowledgment number (32 bits)  |
| If the ACK flag is set then the value of this field is the next sequence number that the sender      |
| of the ACK is expecting. This acknowledges receipt of all prior bytes (if any). The first ACK        |
| sent by each end acknowledges the other end's initial sequence number itself, but no data.           |
| Data offset (4   |

| bits)   |
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| Specifies the size of the TCP header in 32-bit words. The minimum size header is 5 words and  |
| the maximum is 15 words thus giving the minimum size of 20 bytes and maximum of 60 bytes,     |
| allowing for up to 40 bytes of options in the header. This field gets its name from the fact  |
| that it is also the offset from the start of the TCP segment to the actual data.              |
| Reserved (3 bits)   |
| For future use and should be set to zero.   |
| Flags (9 bits)  |
| Contains 9 1-bit flags (control bits) as follows:   |
| · NS (1 bit): ECN-nonce - concealment protection  |
| · CWR (1 bit): Congestion window reduced (CWR) flag is set by the sending host to indicate    |
| that it received a TCP segment with the ECE flag set and responded in a congestion control    |
| mechanism.  |
| · ECE (1 bit): ECN-Echo has a dual role, depending on the value of the SYN flag. It indicates |
| · If the SYN flag is set (1), that the TCP peer is ECN capable.                               |
| · If the SYN flag is clear (0), that a packet with Congestion Experienced flag set (ECN=11)   |
| in the IP header was received during normal transmission. This serves as an indication of     |
| network congestion (or impending congestion) to the TCP sender.                               |
| · URG (1 bit): Indicates that the Urgent pointer field is significant                         |
| · ACK (1 bit): Indicates that the Acknowledgment field is significant. All packets after the  |
| initial SYN packet sent by the client should have this flag set.                              |
| · PSH (1 bit): Push function. Asks to push the buffered data to the receiving application.    |
| · RST (1 bit): Reset the connection   |
| · SYN (1 bit): Synchronize sequence numbers. Only the first packet sent from each end should  |
| have this flag set. Some other flags and fields change meaning based on this flag, and some   |
| are only valid when it is set, and others when it is clear.                                   |
| · FIN (1 bit): Last packet from sender  |
| Window size (16   |

| / w 1  |
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| bits)  |
| The size of the receive window, which specifies the number of window size units that the         |
| sender of this segment is currently willing to receive.  |
| Checksum (16 bits)   |
| The 16-bit checksum field is used for error-checking of the TCP header, the payload and an       |
| IP pseudo-header. The pseudo-header consists of the source IP address, the destination IP        |
| address, the protocol number for the TCP protocol (6) and the length of the TCP headers and      |
| payload (in bytes).  |
| Urgent pointer (16 bits)   |
| If the URG flag is set, then this 16-bit field is an offset from the sequence number indicating  |
| the last urgent data byte.   |
| Options (Variable 0–320 bits, in units of 32 bits)   |
| TCP provides the concept of a connection. A process creates a TCP socket by calling the          |
| socket() function with the parameters PF_INET or PF_INET6 and SOCK_STREAM.                       |
| Server   |
| Setting up a simple TCP server involves the following steps:                                     |
| 1. Creating a TCP socket, with a call to socket().   |
| 2. Binding the socket to the listen port, with a call to bind(). Before calling bind(),          |
| sockaddr_in structure must be declared, and it and the sin_family (AF_INET or AF_INET6)          |
| must be cleared, and fill its sin_port (the listening port, in network byte order) fields.       |
| Converting a short int to network byte order can be done by calling the function htons()         |
| (host to network short).   |
| 3. Preparing the socket to listen for connections (making it a listening socket), with a call to |
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| listen().   |
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| 4. Accepting incoming connections, via a call to accept(). This blocks until an incoming  |
| connection is received, and then returns a socket descriptor for the accepted connection. The   |
| initial descriptor remains a listening descriptor, and accept() can be called again at any time with this socket, until it is closed. |
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| 5. Communicating with the remote host, which can be done through send() and recv().   |
| 6. Eventually closing each socket that was opened, once it is no longer needed, using close().  |
| Note that if there were any calls to fork(), each process must close the sockets it knew about  |
| (the kernel keeps track of how many processes have a descriptor open), and two processes  |
| should not use the same socket at once.   |
| Client  |
| Setting up a TCP client involves the following steps:   |
| 1. Creating a TCP socket, with a call to socket().  |
| 2. Connecting to the server with the use of connect, passing a sockaddr_in structure with the   |
| sin_family set to AF_INET or AF_INET6, sin_port set to the port the endpoint is listening   |
| (in network byte order), and sin_addr set to the IPv4 or IPv6address of the listening server  |
| (also in network byte order.)   |
| 3. Communicating with the server by send()ing and recv()ing.Terminating the connection and  |
| cleaning up with a call to close(). Again, if there were any calls to fork(), each process must                                       |
| close() the socket.   |
| Conclusion:   |
| Thus we have successfully implemented the socket programming for TCP using C.   |
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