

# CS406: Compilers

HA-1

190010008

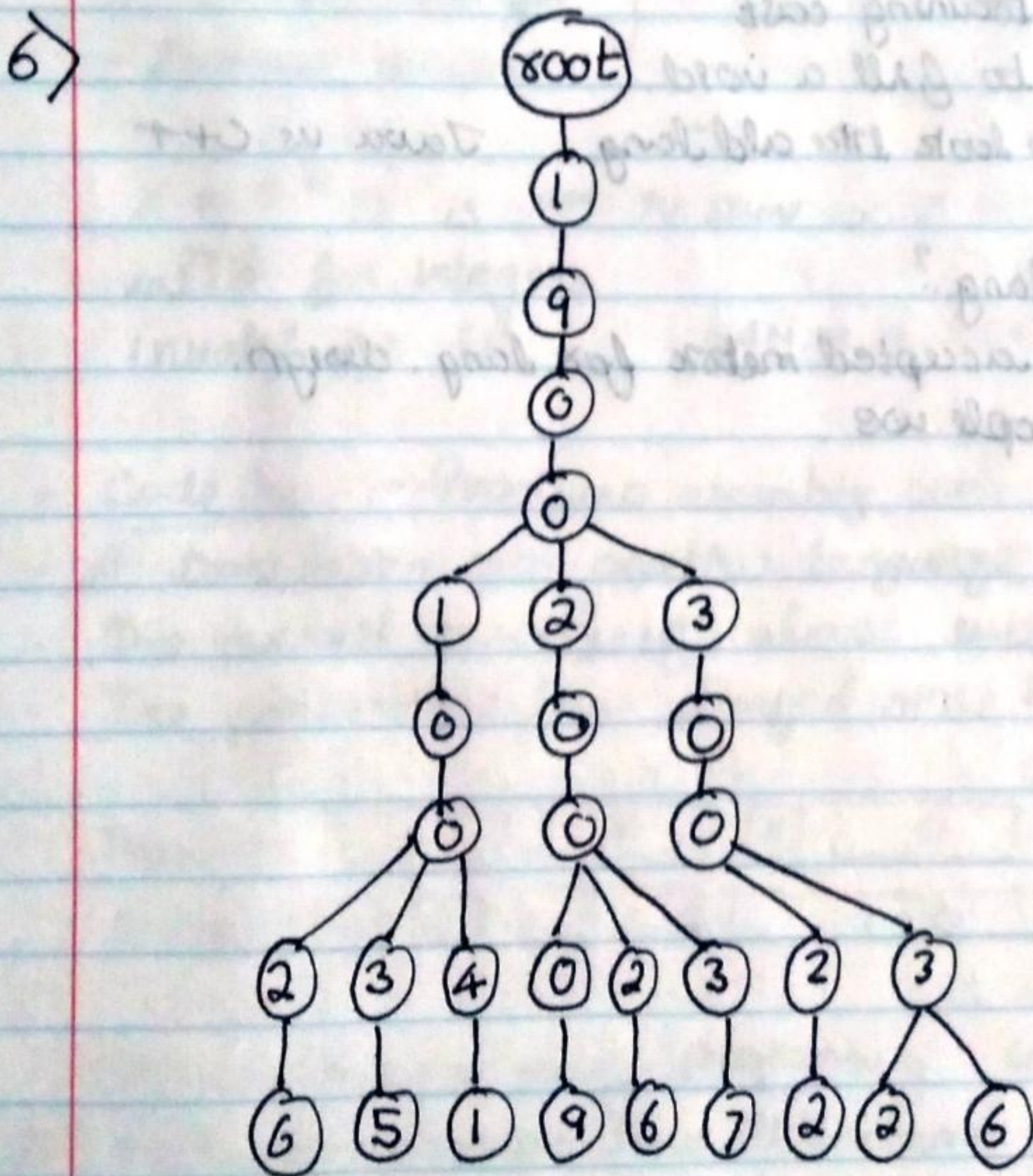
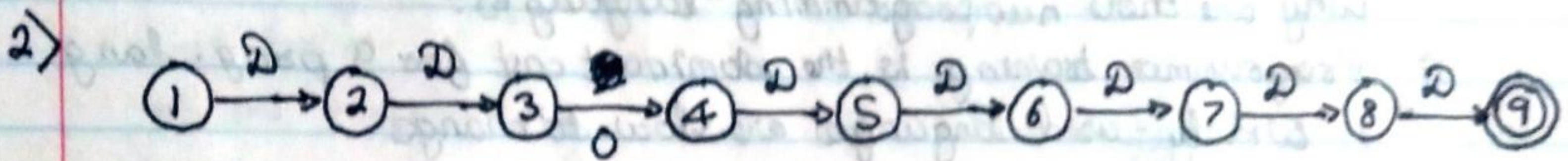
1) Regular Expression that matches valid roll number (aabbcc) ⇒

$[0-9][0-9]0[0-9][0-9][0-9][0-9][0-9]$

OR

~~DDDDDDDD~~  
DD0DDDD

where  $D = [0-9]$



CSE

190010026

190010035

190010041

ECE

190020009

190020026

190020037

MMAE

190030022

190030032

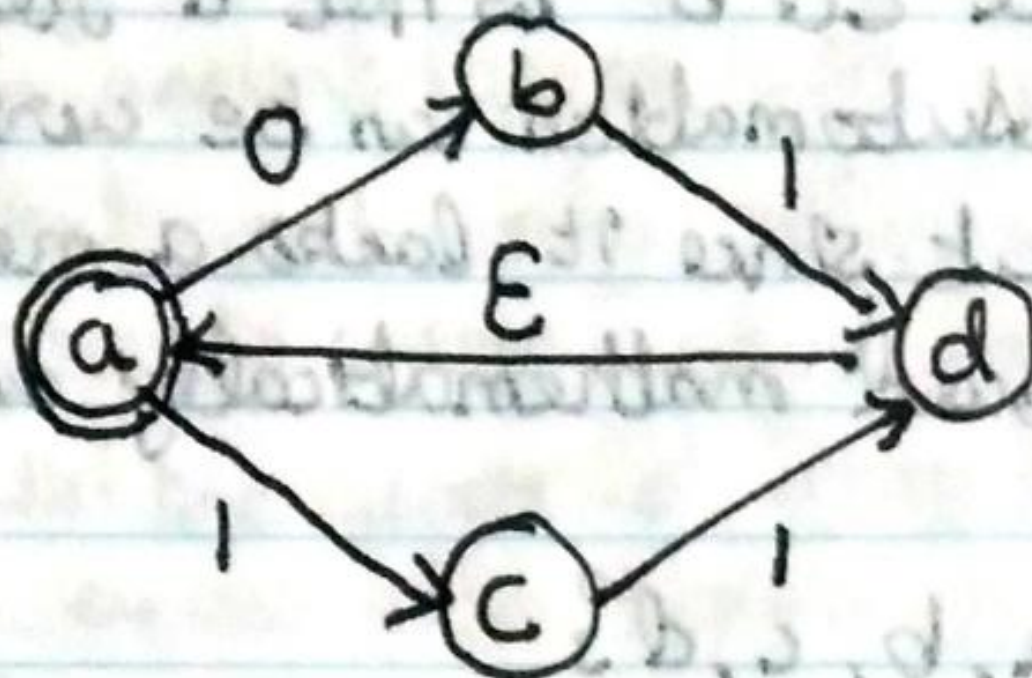
190030036



3) Regular Expression  $\Rightarrow (((01) + 1(11) +))^*$

Strings belonging to the above regex =  $\epsilon, 01, 11, 0111, 0101, 1111, 010111, \text{etc}$

NFA  $\Rightarrow$

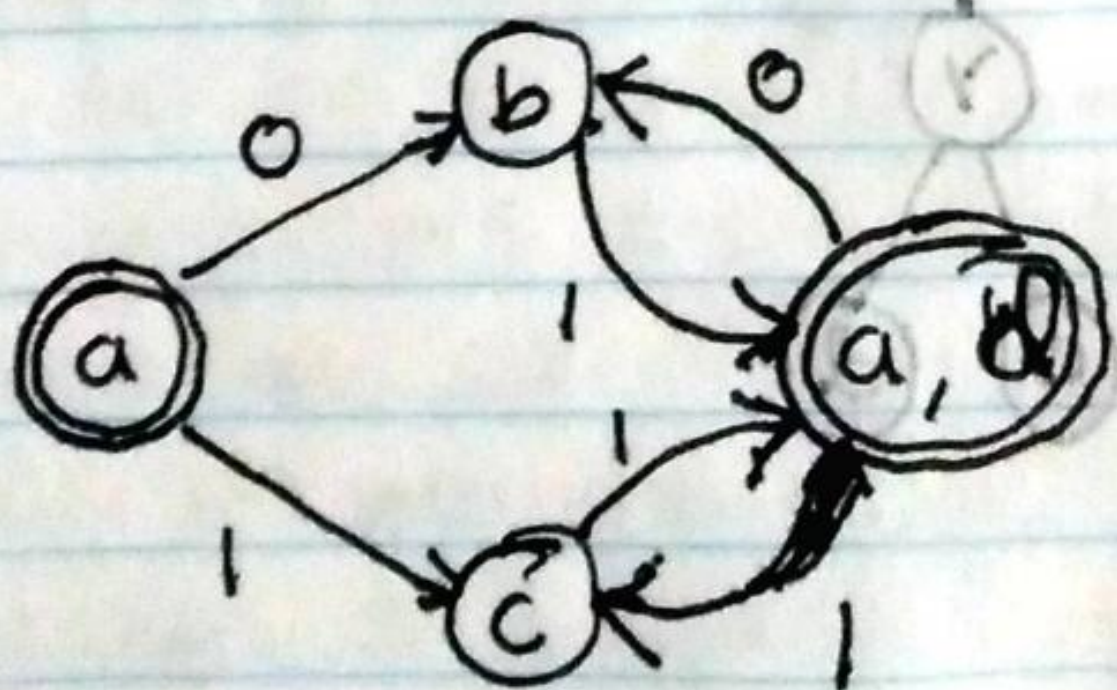


a is the initial as well as the final state

4) Transition Table  $\Rightarrow$

State \ Character	$\epsilon$	0	1
a	a	b	d
b	b	-	d
c	c	-	d
d	a, d	-	-
a, d	a, d	b	c

Equivalent DFA  $\Rightarrow$



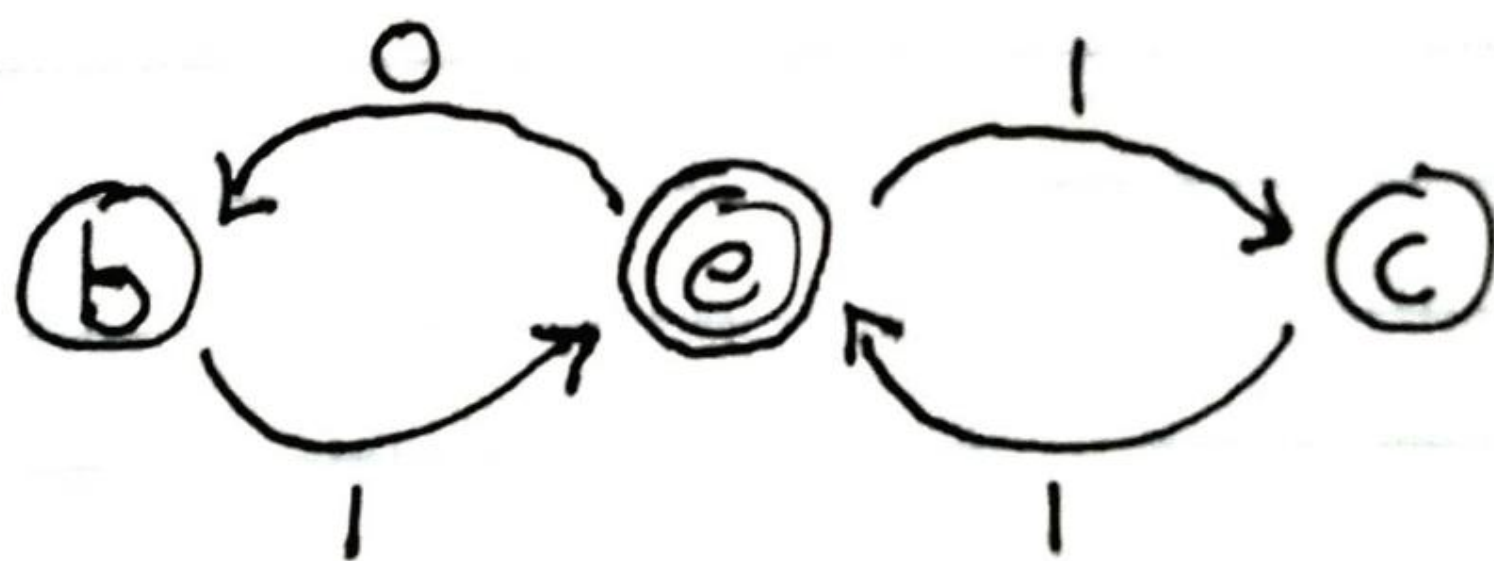
a is the initial state & (a) and (a,d) are the final state.



- 5) In order to convert DFA to reduced DFA, we look into the transition table to find which states could be merged together. We see that state (a) & (a,d) could be merged.  
 State Transition Table after merging -

State \ Character	ε	0	1
a-a,d-M	a-a,d-M	b	c
b	b	-	a-a,d-M
c	c	-	a-a,d-M

Lets rename the state a-a,d-M as state e



e is the initial as well as the final state



7) Strings with even number of characters having  $n$  number of a's followed by cd followed by  $n$  number of b's where  $n \geq 0$ .

8) The language  $a^n cd b^n$  is not a regular expression since no Finite State Automaton can be written that validates & accepts this input since it lacks a memory component. This can also be proved mathematically using the pumping lemma.

9) Terminals:  $\{a, b, c, d\}$   
Non-Terminals:  $\{S, X, Y\}$

Step	Production Used
$S \rightarrow X$	$S \rightarrow X$
$X \rightarrow axb$	$X \rightarrow axb$
$axb \rightarrow aaxbb$	$X \rightarrow axb$
$aaxbb \rightarrow aaybb$	$X \rightarrow Y$
$aaybb \rightarrow aacd bbb$	$Y \rightarrow cd$

Parse Tree:

