



Academic Year (2021-22)

Year: 2 Semester: IV

Program: B. Tech. (Computer Engg.)

Subject: Formal Language and Automata Theory

Date:

Max. Marks: 75

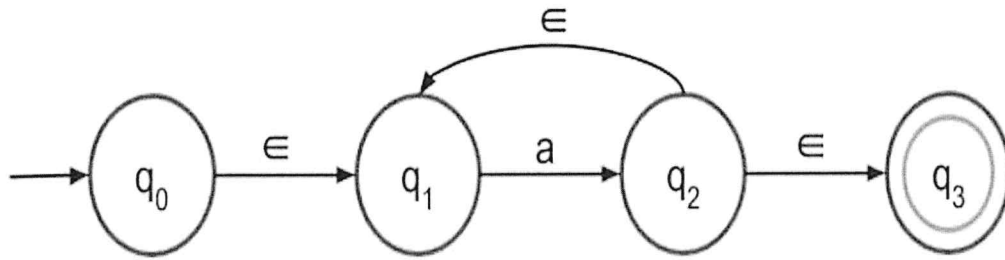
Time: 10: 30 am to 1:30 pm

Duration: 3 Hours

RE . EXAMINATION

ANSWER KEY

| Question No. | | Max. Marks |
|--------------|---|------------|
| Q1 (a) | <p>Step 1: STF 2marks Step 2: MTF 2marks Step3: Transition diagram 1 mark</p> <p><i>state transition diagram</i></p> | [05] |
| Q1 (b) | <p>Step1: Definition 2marks Step2: Transition Table 2marks Step3: Implementation 4marks Step4: verify string 2marks OR Step1: RE to NFA 3marks Step2: NFA to DFA 4marks Step3: DFA to Min DFA 4marks</p> | [10] |



Step1: Simplify grammar 4marks
 Step2: Convert simplified grammar into CNF 4marks

Q2 (a)

[8]

$A \rightarrow aBa \mid bBa$
 $B \rightarrow aB \mid bB \mid \epsilon$

soln...

Step 1: Remove ϵ -prodn

$A \rightarrow aBa \mid bBa \mid aa \mid ba$
 $B \rightarrow aB \mid bB \mid a \mid b$

Step 2: CNF format

$A \rightarrow a \text{ or } A \rightarrow BC$

Step 3: $\begin{cases} X \rightarrow a \\ Y \rightarrow b \end{cases}$

$A \rightarrow XB \mid YB \mid XX \mid YX$
 $B \rightarrow XB \mid YB \mid a \mid b$

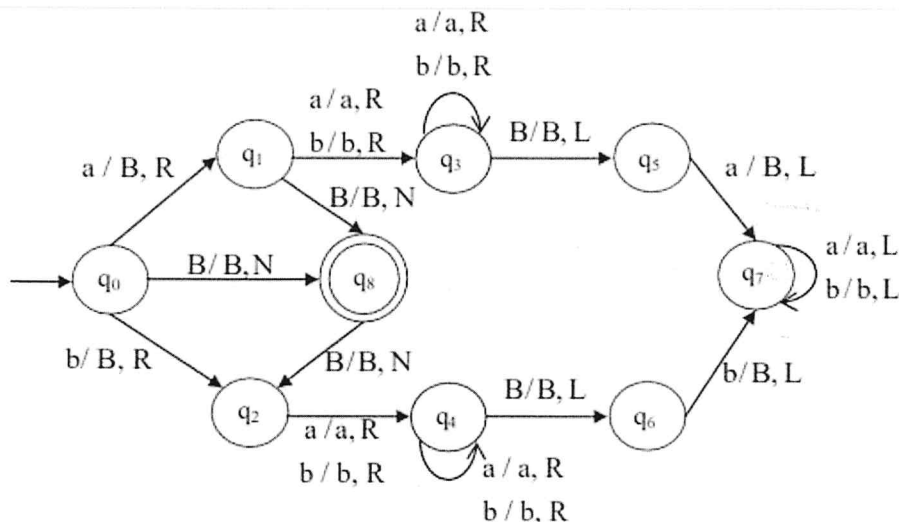
Let $\begin{cases} U \rightarrow XB \\ V \rightarrow YB \end{cases}$

Step 4:

$A \rightarrow UX \mid VX \mid XX \mid YX$
 $B \rightarrow XB \mid YB \mid a \mid b$
 $X \rightarrow a$
 $Y \rightarrow b$
 $U \rightarrow XB$
 $V \rightarrow YB$



| | | |
|--------|---|-----|
| | | |
| Q2 (b) | <p style="text-align: center;">OR</p> <p>Step1: Leftmost Derivation 2.5marks Step2: Rightmost Derivation 2.5marks Step3: Parse tree 2marks</p> <p>Leftmost derivation:</p> $S \Rightarrow S + S$ $S \Rightarrow SS + S \text{ (as } S \Rightarrow SS)$ $S \Rightarrow a^*S + S \text{ (as } S \Rightarrow a)$ $S \Rightarrow a^*b + S \text{ (as } S \Rightarrow b)$ $S \Rightarrow ab + SS \text{ (as } S \Rightarrow S^*S)$ $S \Rightarrow ab + aS \text{ (as } S \Rightarrow a)$ $S \Rightarrow ab + ab \text{ (as } S \Rightarrow b)$ <p>Rightmost derivation:</p> $S \Rightarrow S + S$ $S \Rightarrow S + SS \text{ (as } S \Rightarrow SS)$ $S \Rightarrow S + S^*b \text{ (as } S \Rightarrow b)$ $S \Rightarrow S + a^*b \text{ (as } S \Rightarrow a)$ $S \Rightarrow SS + ab \text{ (as } S \Rightarrow S^*S)$ $S \Rightarrow Sb + ab \text{ (as } S \Rightarrow b)$ $S \Rightarrow ab + ab \text{ (as } S \Rightarrow a)$ | [7] |
| Q3 (a) | <p>Step1: Definition of PDA 2marks Step2: Logic 2marks Step3: Implementation 3marks</p> <p>OR</p> <p>Step1: Definition of PDA 2marks Step2: Logic 2marks Step3: Implementation 3marks</p> | [7] |
| Q3 (b) | <p>Step1: Definition of TM 2marks Step2: Logic 2mark Step3: Implementation 3marks Step4: Example 1Mark</p> | [8] |



Step1: Let L be the regular language. 1Mark
 Step2: Select a string Z as per given condition. 3 Marks
 Step3: Get required equation using pumping lemma. 3Marks
 Step4: Prove that our assumption is wrong. 3Marks

$$L_> = \{a^i b^j : i > j\}$$

$L_>$ is not regular.

- Fix an arbitrary pumping length $n > 0$.
- Choose a proper string s in $L_>$.
- $s = a^{n+1} b^n \in L_>$.
- Consider all possible splittings of s in x, y, z with the desired properties: $y = a^m, 1 \leq m \leq n$.
- $xz = a^{n+1-m} b^n \notin L_>$.
- So $L_>$ is not regular!

The strings which end with 00 or 11. Let us find out L_1 and L_2 . Hence $R = [(00+11)(0+1)^*] + [(0+1)^*(00+11)]$

NFA diagram

Diagram



| | | |
|---------------|--|--|
| | <div data-bbox="268 286 1038 481"></div> <p>And explanation.</p> <p>OR</p> <ul style="list-style-type: none">• A Turing machine is said to be universal Turing machine if it can accept:<ul style="list-style-type: none">◦ The input data, and◦ An algorithm (description) for computing. | <p>[3]</p> <p>(5)</p> <p>(10)</p> |
| <p>Q5 (b)</p> | <ol style="list-style-type: none">1. Type 0 is known as unrestricted grammar.2. Type 1 is known as context-sensitive grammar.3. Type 2 is known as a context-free grammar.4. Type 3 Regular Grammar. <div data-bbox="263 1030 1204 1579"></div> <p>Explanation with grammar as examples</p> | <p>[2]</p> <p>[1]</p> <p>[2]</p> |

