

unit - 3

Dynamic Programming

→ All pair shortest path algorithm

- Floyd Warshall's Algorithm

$$G(V, E)$$

$V \rightarrow$ set of vertices

$E \rightarrow$ set of edges

$m \rightarrow$ no. of vertices (K)

$i \rightarrow$ source node or outgoing edge

$j \rightarrow$ destination or incoming edge

If self or parallel loop is given then we don't have to solve it

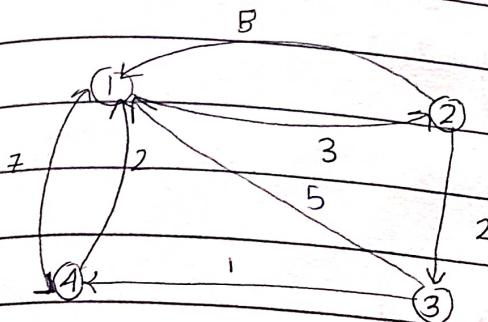
$$D^0 = W$$

$$D_{ij}^k = \begin{cases} W_{ij} & \dots \text{for } k=0 \\ \min \left[D_{ij}^{(k-1)}, D_{ik}^{(k-1)} + D_{kj}^{(k-1)} \right] & \dots \text{for } k \geq 1 \end{cases}$$

(1) (2)

$$\text{Time complexity} = O(n^3)$$

(a) Find out the shortest path in the diagram



$$n = \text{no. of node} = 4$$

$$K \rightarrow 0 \text{ to } 4$$

$$D^K = D^0$$

	1	2	3	4
1	0	3	∞	7
2	8	0	2	∞
3	5	∞	0	1
4	2	∞	0	0

For $K = 1$

$$D^1 =$$

	1	2	3	4
1	0	3	∞	7
2	8	0	2	15
3	5	∞	0	1
4	2	5	∞	0

negate the 5th row & 1st col in D^0

$$D_{ij}^k = \min(D_{23}^0, D_{21}^0 + D_{13}^0)$$

$$= \min(2, 8 + \infty)$$

$$D_{23}^0 = 2$$

$$D_{ij}^k = \min (D_{ij}^{K=1}, D_{iH}^0 + D_{kj}^{K=1})$$

$$\begin{aligned} D_{21}^0 &= \min (D_{2A}^0, D_{21}^0 + D_{11}^0) \\ &= \min (\infty, 8+7) \\ &= 15 \end{aligned}$$

$$\begin{aligned} D_{32}^1 &= \min (D_{32}^0, D_{31}^0 + D_{12}^0) \\ &= \min (\infty, 5+3) \\ &= 8 \end{aligned}$$

$$\begin{aligned} D_{33}^1 &= \min (D_{33}^0, D_{31}^0 + D_{13}^0) \\ &= \min (\infty, 5+\infty) \\ &= \infty \end{aligned}$$

$$\begin{aligned} D_{34}^1 &= \min (D_{34}^0, D_{31}^0 + D_{14}^0) \\ &= \min (\infty, 5+7) \\ &= 1 \end{aligned}$$

$$\begin{aligned} D_{42}^1 &= \min (D_{42}^0, D_{41}^0 + D_{12}^0) \\ &= \min (\infty, 2+3) \\ &= 5 \end{aligned}$$

$$\begin{aligned} D_{43}^1 &= \min (D_{43}^0, D_{41}^0 + D_{13}^0) \\ &= \min (\infty, 2+\infty) \\ &= \infty \end{aligned}$$

For K=2

$$D^2 = \begin{array}{|c|c|c|c|} \hline 0 & 3 & 5 & 7 \\ \hline 8 & 0 & 2 & 15 \\ \hline 5 & 8 & 0 & 1 \\ \hline 2 & 5 & 7 & 0 \\ \hline \end{array}$$

$$\begin{aligned} D_{12}^2 &= \min(D_{12}^1, D_{12}^1 + D_{22}^1) \\ &= \min(3, 3+0) \\ &= 3 \end{aligned}$$

$$\begin{aligned} D_{13}^2 &= \min(D_{13}^1, D_{12}^1 + D_{23}^1) \\ &= \min(\infty, 3+2) \\ &= 5 \end{aligned}$$

$$\begin{aligned} D_{14}^2 &= \min(D_{14}^1, D_{12}^1 + D_{24}^1) \\ &= \min(7, 3+15) \\ &= 7 \end{aligned}$$

$$\begin{aligned} D_{22}^2 &= \min(D_{22}^1, D_{32}^1 + D_{22}^1) \\ &= \min(8, 8+0) \\ &= 8 \end{aligned}$$

$$\begin{aligned} D_{34}^2 &= \min(D_{34}^1, D_{32}^1 + D_{24}^1) \\ &= \min(1, 8+15) \\ &= 1 \end{aligned}$$

Now $K = 3$

$$D^3 = \begin{array}{|c|c|c|c|} \hline 0 & 3 & 5 & 6 \\ \hline 7 & 0 & 2 & 3 \\ \hline 5 & 8 & 0 & 1 \\ \hline 2 & 5 & 7 & 0 \\ \hline \end{array}$$

$$D_{12}^3 = \min(D_{12}^2, D_{13}^2 + D_{32}^2)$$

$$= \min(3, 5+8)$$

$$= 3$$

$$D_{14}^3 = \min(D_{14}^2, D_{13}^2 + D_{34}^2)$$

$$= \min(7, 5+1)$$

$$= 6$$

$$D_{21}^3 = \min(D_{21}^2, D_{23}^2 + D_{31}^2)$$

$$= \min(8, 2+5)$$

$$= 7$$

$$D_{24}^3 = \min(D_{24}^2, D_{23}^2 + D_{34}^2)$$

$$= \min(15, 2+1)$$

$$= 3$$

$$D_{41}^3 = \min(D_{41}^2, D_{43}^2 + D_{31}^2)$$

$$= \min(2, 7+5)$$

$$= 2$$

$$D_{42}^3 = \min(D_{42}^2, D_{43}^2 + D_{32}^2)$$

$$= \min(5, 7+8)$$

$$= 5$$

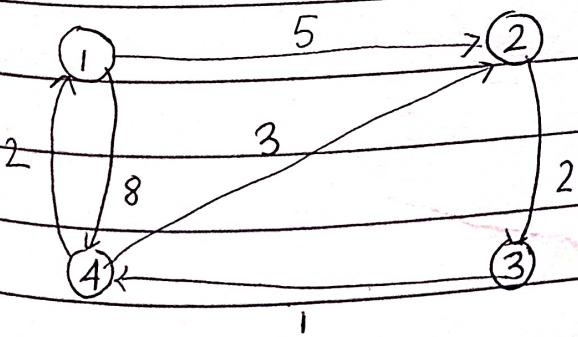
Now $K=4$

$D^4 =$	0	3	5	6
	5	0	2	3
	3	6	0	1
	2	5	7	0

$$D_{12}^4 = \min(D_{12}^3, D_{14}^3 + D_{42}^3)$$

$$= \min(3, 6 + 5)$$

$$= 3$$



Now For K=0

	1	2	3	4	
D ⁰ =1	0	5	∞	8	
2	∞	0	2	∞	
3	∞	∞	0	1	
4	2	3	∞	0	

Now for K=1

$$D_{ij}^K = \min(D_{ij}^{K-1}, D_{ik}^{K-1} + D_{kj}^{K-1})$$

$$\begin{aligned} D_{13}^1 &= \min(D_{13}^0, D_{11}^0 + D_{13}^0) \\ &= \min(\infty, 0 + \infty) \\ &= \infty \end{aligned}$$

0

D' =	0	5	∞	8	
	∞	0	2	∞	
	∞	∞	0	1	
	2	3	∞	0	

$$\begin{aligned} D_{23}' &= \min(D_{23}^0, D_{21}^0 + D_{13}^0) \\ &= \min(2, \infty + \infty) \\ &= 2 \end{aligned}$$

$$\begin{aligned} D_{24}' &= \min(D_{24}^0, D_{21}^0 + D_{14}^0) \\ &= \min(\infty, \infty + 8) \\ &= \infty \end{aligned}$$

$$D'_{32} = \min(D^0_{32} + D^0_{31} + D^0_{12})$$

$$= \min(\infty, \infty + 5)$$

$$= \infty$$

$$D'_{34} = \min(D^0_{34}, D^0_{31} + D^0_{14})$$

$$= \min(1, \infty + 8)$$

$$= 1$$

$$D'_{42} = \min(D^0_{42}, D^0_{41} + D^0_{12})$$

$$= \min(3, 2 + 5)$$

$$= 3$$

$$D'_{43} = \min(D^0_{43}, D^0_{41} + D^0_{13})$$

$$= \min(\infty, 2 + \infty)$$

$$= \infty$$

Now For K=2

	0	5	7	8
D ² =	∞	0	2	∞
	0	∞	0	1
	2	3	5	0

$$D^2_{13} = \min(D'_{13}, D'_{12} + D'_{23})$$

$$= \min(\infty, 5 + 2)$$

$$= 7$$

$$D^2_{14} = \min(D'_{14}, D'_{12} + D'_{24})$$

$$= \min(8, 5 + \infty)$$

$$= 8$$

$$D_{31}^2 = \min(D_{31}', D_{32}' + D_{21}')$$

$$= \min(\infty, \infty + \infty)$$

$$= \infty$$

$$D_{31}^2 = \min(D_{31}', D_{32}' + D_{21}')$$

$$= \min(1, \infty + \infty)$$

$$= 1$$

$$D_{41}^2 = \min(D_{41}', D_{42}' + D_{21}')$$

$$= \min(2, 3 + \infty)$$

$$= 2$$

$$D_{43}^2 = \min(D_{43}', D_{42}' + D_{23}')$$

$$= \min(\infty, 3 + 2)$$

$$= 5$$

Now For K=3

	0	5	7	8			
D_{12}^3	∞	0	2	3			
	∞	∞	0	1			
	2	3	5	0			

$$D_{12}^3 = \min(D_{12}^2, D_{13}^2 + D_{32}^2)$$

$$= \min(5, 7 + \infty)$$

$$= 5$$

$$D_{14}^3 = \min(D_{14}^2, D_{13}^2 + D_{34}^2)$$

$$= \min(8, 7 + 1)$$

$$= 8$$

$$D_{21}^3 = \min(D_{21}^2, D_{23}^2 + D_{31}^2)$$

$$= \min(\infty, 2 + \infty)$$

$$= \infty$$

$$D_{24}^3 = \min(D_{24}^2, D_{23}^2 + D_{34}^2)$$

$$= \min(\infty, 2 + 1)$$

$$= 3$$

$$D_{41}^3 = \min(D_{41}^2, D_{43}^2 + D_{31}^2)$$

$$= \min(2, 5 + \infty)$$

$$= 2$$

$$D_{12}^3 = \min(D_{42}^2, D_{43}^2 + D_{32}^2)$$

$$= \min(3, 5 + \infty)$$

$$= 3$$

Now For $K=4$

$$D^4 = \begin{array}{|c|c|c|c|} \hline & 0 & 5 & 7 & 8 \\ \hline 0 & & & & \\ \hline 5 & & 0 & 2 & 3 \\ \hline 3 & 4 & & 0 & 1 \\ \hline 2 & 3 & 5 & & 0 \\ \hline \end{array}$$

$$D_{12}^4 = \min(D_{12}^3, D_{14}^3 + D_{42}^3)$$

$$= \min(5, 8 + 3) = 5$$

$$D_{13}^4 = \min(D_{13}^3, D_{14}^3 + D_{43}^3)$$

$$= \min(7, 8 + 5) = 7$$

$$D_{21}^4 = \min(D_{21}^3, D_{24}^3 + D_{41}^3)$$

$$= \min(\infty, 3 + 2) = 5$$

$$D_{23}^4 = \min(D_{23}^3, D_{24}^3 + D_{43}^3)$$

$$= \min(2, 3+5)$$

$$= 2$$

$$D_{31}^4 = \min(D_{31}^3, D_{34}^3 + D_{41}^3)$$

$$= \min(\infty, 1+2)$$

$$= 3$$

$$D_{32}^4 = \min(D_{32}^3, D_{34}^3 + D_{42}^3)$$

$$= \min(\infty, 1+3)$$

$$= 4$$

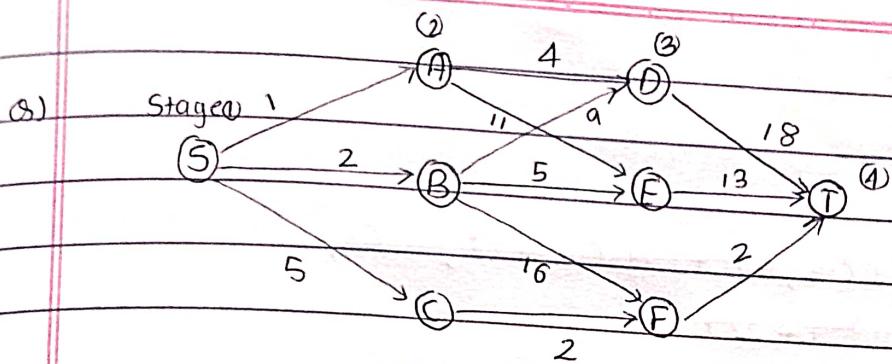
* Multistage Graph

\rightarrow

$$MS(s_i, v_j) = \min_{\substack{\text{Stage} \\ \text{Vertex}}} (c(v_j, k) + MS(s_{j+1}, k))$$

next stage

$\forall k \in S_{i+1}; (v_j, k) \in E \dots \text{edge}$



$$\rightarrow \text{MSC}(S_i, v_j) = \min(c(v_j, K) + \text{MS}(S_{j+1}, K))$$

Stage 1

$$\text{MS}(A, T) = \text{dest}^n \text{ node} = 0$$

\rightarrow Stage 3

$$\text{MS}(3, D) = \min(C(D, T) + \text{MS}(4, T)) = \min(18 + 0) = \min\{18\}$$

$$\text{MS}(3, F) = \min(C(F, T) + \text{MS}(4, T)) = \min(13 + 0) = 13$$

$$\text{MS}(3, F) = \min(C(F, T) + \text{MS}(4, T)) = \min(2 + 0) = 2$$

\rightarrow Stage 2

$$\begin{aligned} \text{MS}(2, A) &= \min\{C(A, D) + \text{MS}(3, D); &= \min\{C(A+18), (11+13)\} \\ \text{MS}(2, A) &\quad C(A, E) + \text{MS}(3, E)\} &= 22 \end{aligned}$$

$$\begin{aligned} \text{MS}(2, B) &= \min\{C(B, D) + \text{MS}(3, D); &= \min\{(9+18), (5+13), (16+2)\} \\ \text{MS}(2, B) &\quad C(B, E) + \text{MS}(3, E); &= 18 \\ \text{MS}(2, B) &\quad C(B, F) + \text{MS}(3, F)\} \end{aligned}$$

$$\begin{aligned} \text{MS}(2, C) &= \min(C(C, F) + \text{MS}(3, F)) = \min\{2 + 2\} \\ &= 4 \end{aligned}$$

Stage 1 :

$$\begin{aligned} MS(1, S) &= \min \{ CC(S, A) + MSC(2, A) \}; CC(S, \\ &= \min \{ 1 + 22 \} \\ &= 23 \end{aligned}$$

MS(1,)

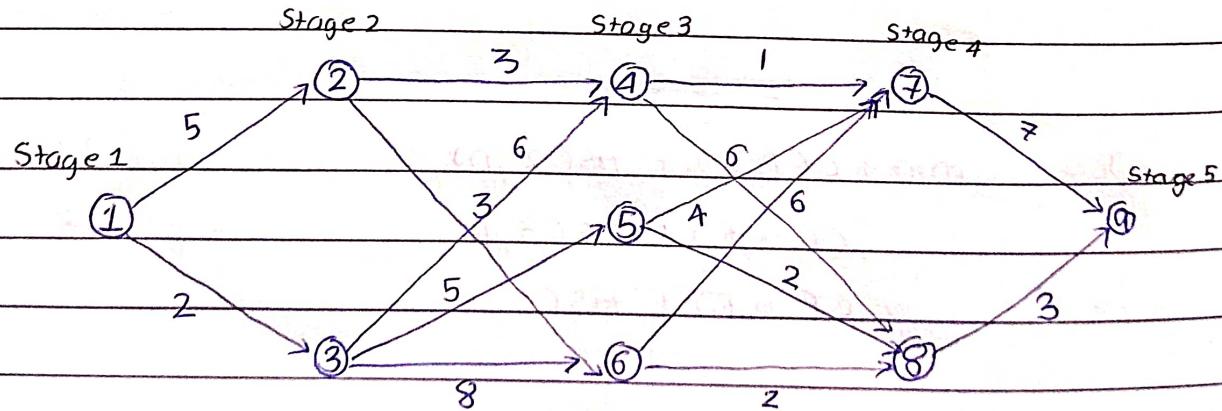
$$\begin{aligned} MS(1, S) &= \min \{ CC(S, A) + MS(2, A); CC(S, B) + MS(2, B); \\ &\quad CC(S, C) + MSC(2, C) \} \\ &= \min \{ (1+22); (2+18); (5+4) \} \\ &= 9 \end{aligned}$$

So shortest distance betⁿ S to T is 9

& Root is

$$S \xrightarrow{5} C \xrightarrow{2} F \xrightarrow{2} T$$

(b) Calculate shortest path & distance from node 1 to 9



$$g(i, \emptyset) = C_{i1}$$

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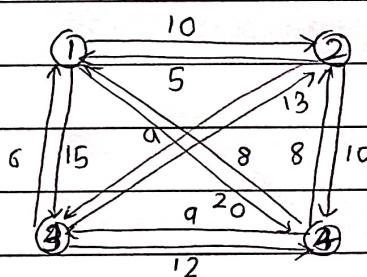
* Travelling Salesman problem

→ Self loop & Parallel edges are not allowed

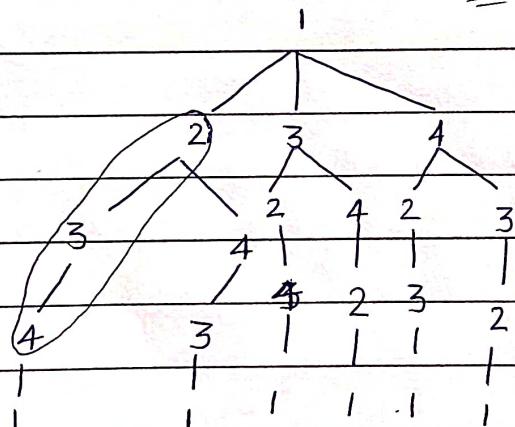
→ Salesman have to start with vertex 1, visit all the nodes and return back to vertex 1

→ For Brute force approach Time complexity, is

$$O(n!) = O(n^n) = \text{Large T.C}$$



Brute force
= = =



$$\text{Formula} = G(i, S) = \min_j \{ C_{ij} + g(j, S - \{ j \}) \}$$

↓
vertices $j \in S$

$S \rightarrow$ sets of states or vertices

→ There are almost $2^n \cdot n$ subproblems & each one takes linear time to solve

$$\therefore \text{TC is } \Theta(n!) \text{ TSP is } O(2^n \cdot n^2)$$

	1	2	3	4
1	0	10	15	20
2	5	0	9	10
3	6	13	0	12
4	8	8	9	0

→ If edge is not available consider as ∞

take $|S| = 0$, $S = \emptyset$

$$G(2, \emptyset) = c_{12} = c_{21} = 5$$

$$G(3, \emptyset) = c_{13} = c_{31} = 6$$

$$G(4, \emptyset) = c_{14} = c_{41} = 8$$

take $|S| = 1$

$$G(i, S) = \min (c_{ij} + g(j, S - \{j\})) \rightarrow \{3\} - \{3\} = \emptyset$$

$$\begin{aligned} G(2, \{3\}) &= \min (c_{23} + g(3, \emptyset)) \\ &= \min (9 + 6) = 15 \end{aligned}$$

$$\begin{aligned} G(2, \{4\}) &= \min (c_{24} + g(4, \emptyset)) \\ &= \min (10 + 8) = 18 \end{aligned}$$

$$\begin{aligned} G(3, \{4\}) &= \min (c_{34} + g(4, \emptyset)) \\ &= \min (12 + 8) \\ &= 20 \end{aligned}$$

$$\begin{aligned} G(3, \{2\}) &= \min (c_{32} + g(2, \emptyset)) \\ &= \min (13 + 5) \\ &= 18 \end{aligned}$$

$$G(1, \{2\}) = \min(C_{12} + g(2, \emptyset))$$

$$= \min(8 + \frac{5}{8})$$

$$= 16.13$$

$$G(4, \{3\}) = \min(C_{13} + g(3, \emptyset))$$

$$= \min(9 + 6)$$

$$= 15$$

for $|S|=2$

$$G(2, \{3, 4\}) = \min \left\{ C_{23} + G(3, \{4\}), \right.$$

$$\left. C_{24} + G(4, \{3\}) \right\}$$

$$= \min \left\{ 9 + 20, \right.$$

$$\left. 10 + 15 \right\} = \min \{29, 25\} = 25$$

$$G(3, \{2, 4\}) = \min \left\{ C_{32} + G(2, \{4\}), \right.$$

$$\left. C_{34} + G(4, \{2\}) \right\}$$

$$= \min \left\{ 13 + 18, \right.$$

$$\left. 12 + 13 \right\}$$

$$= \min(31, 25) = 25$$

$$G(4, \{2, 3\}) = \min \left\{ C_{42} + G(2, \{3\}), \right.$$

$$\left. C_{43} + G(3, \{2\}) \right\}$$

$$= \min \left\{ 8 + 15, \right.$$

$$\left. 9 + 18 \right\}$$

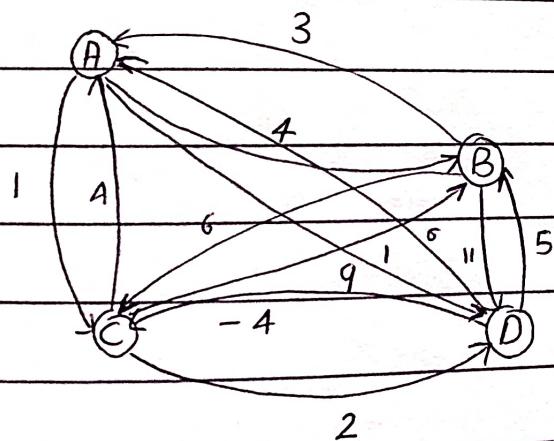
$$= \min \{23, 27\}$$

$$= 23$$

$$|S| = 3 = \{2, 3, 4\}$$

$$\begin{aligned}
 G(1, \{2, 3, 4\}) &= \min \left\{ \begin{array}{l} C_{12} + G(2, \{3, 4\}) \\ C_{13} + G(3, \{2, 4\}) \\ C_{14} = G(4, \{2, 3\}) \end{array} \right\} \\
 &= \min \left\{ \begin{array}{l} 10 + 25, \\ 15 + 25, \\ 20 + 23 \end{array} \right\} \\
 &= 35 \rightarrow \text{SP cost}
 \end{aligned}$$

Shortest path - 1 - 2 - 4 - 3 - 1



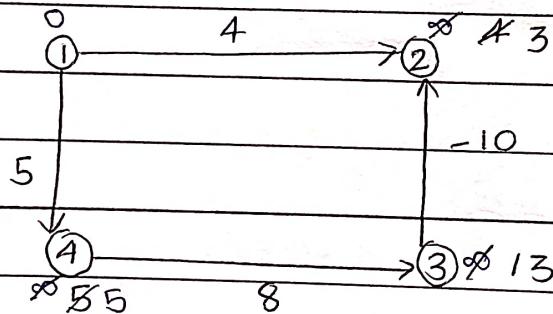
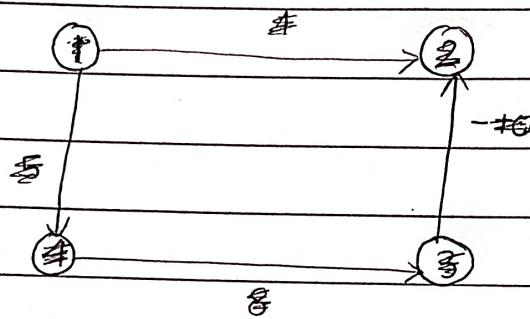
source \rightarrow inner degree 0
dest \rightarrow out " 0

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* Single source shortest path Bellman Ford Algorithm

\rightarrow We can't solve the -ve weight cycle

\rightarrow N no. of vertices No of Iterⁿ is n-1



1-2

1-4

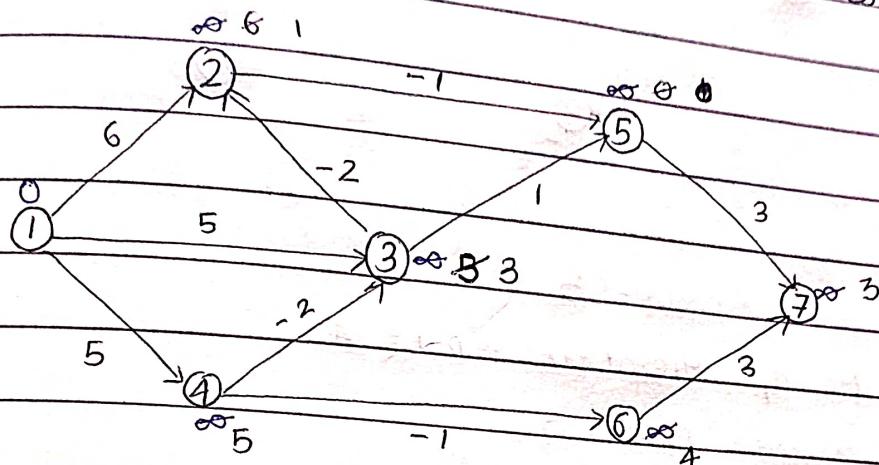
4-3

3-2

Ver gter	1	2	3	4	
Initial cond ⁿ \rightarrow	0	0	∞	∞	∞
1	0	*3	13	5	
2	0	3	13	5	
3	0	3	13	5	

Time complexity : $O(VE)$

Q) Find out the shortest distance from source to desⁿ using SSSPA



SOLⁿ

No. of Nodes = 7

No. of Iteration = $N-1 = 7-1 = 6$

Source node = 1

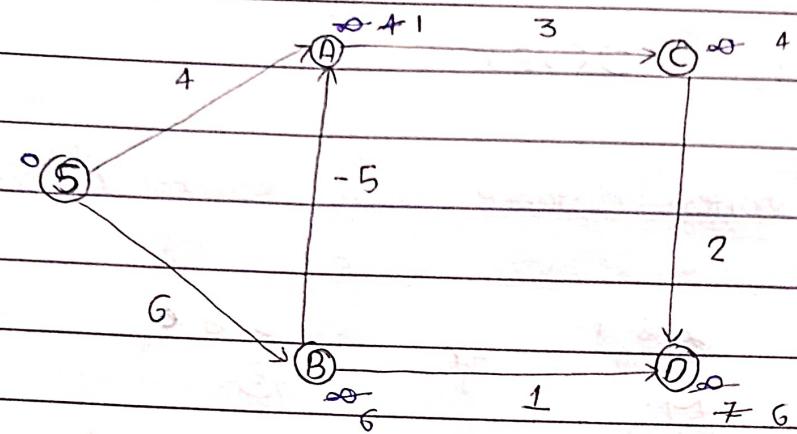
Desⁿ node = 7

Path
 $= 1 \xrightarrow{5} 5 \xrightarrow{-2} 3 \xrightarrow{-2} 2 \xrightarrow{-1} 5 \xrightarrow{3} 7$

Vert Horn	1	2	3	4	5	6	7
0	0	∞	∞	∞	∞	∞	∞
1	0	8 1	5 3	5	0	4	3
2	0	1	3	5	0	4	3
3	0	1	3	5	0	4	3
4	0	1	3	5	0	4	3
5	0	1	3	5	0	4	3
6	0	1	3	5	0	4	3

1-2	✓	✓	3-2	✓	✓	5-7	/
1-4	-	✓	2-5	✓	/	6-7	/
1-3	-	✓	3-5	✓	/		
4-3	✓	/	4-6	✓	/		

(3)

 $\rightarrow \text{Sol}^n$

No. of nodes = 5

No. of iteration = $n-1 = 4$

Source node = S

Destⁿ node = D

Vertic Gter	S	A	B	C	D
0	0	∞	∞	∞	∞
1	0	1	6	4	6
2	0	1	6	4	6
3	0	1	6	4	6
4	0	1	6	4	6

Edges :

S-A = ✓

S-B = ✓

B-A = ✓

A-C = ✓

B-D = ✓

C-D = ✓

Path:

S $\xrightarrow{4}$ B $\xrightarrow{5}$ A $\xrightarrow{3}$ C $\xrightarrow{2}$ D