

MATH220
DISCRETE MATHEMATICS AND CRYPTOGRAPHY

Tutorial 5

Solutions

1. (a) (i) Consider the powers of each of the elements of \mathbb{Z}_{11}^* .

$$\begin{aligned}1^1 &= 1 \\2^1 &= 2, 2^2 = 4, 2^3 = 8, 2^4 = 5, 2^5 = 10, \dots, 2^{10} = 1 \\3^1 &= 3, 3^2 = 9, 3^3 = 5, 3^4 = 4, 3^5 = 1 \\4^1 &= 4, 4^2 = 5, 4^3 = 9, 4^4 = 3, 4^5 = 1 \\5^1 &= 5, 5^2 = 3, 5^3 = 4, 5^4 = 9, 5^5 = 1 \\6^1 &= 6, 6^2 = 3, 6^3 = 7, 6^4 = 9, 6^5 = 10, \dots, 6^{10} = 1 \\7^1 &= 7, 7^2 = 5, 7^3 = 2, 7^4 = 3, 7^5 = 10, \dots, 7^{10} = 1 \\8^1 &= 8, 8^2 = 9, 8^3 = 6, 8^4 = 4, 8^5 = 10, \dots, 8^{10} = 1 \\9^1 &= 9, 9^2 = 4, 9^3 = 3, 9^4 = 5, 9^5 = 1 \\10^1 &= 10, 10^2 = 1\end{aligned}$$

Therefore the order of 1 is 1, the order of 10 is 2, while 3, 4, 5, 9 each have order 5, and 2, 6, 7, 8 each have order 10.

Note. Recall that the order of any element divides the number of elements of \mathbb{Z}_{11}^* , that is, divides $\phi(11) = 10$. So the possible orders are 1, 2, 5, and 10. It follows that if you have calculated g^1, g^2, g^3, g^4 , and g^5 and still have not obtained the value 1, the order must be 10. Knowing this result saves a bit of work!

- (ii) The generators of \mathbb{Z}_{11}^* are 2, 6, 7, and 8.

(b)

$$\begin{aligned}A &\equiv 2^8 \pmod{11} \\&\equiv 256 \pmod{11} \\&\equiv 3 \pmod{11}\end{aligned}$$

So Alice sends $A = 3$ to Bob.

(c)

$$\begin{aligned}B &\equiv 2^6 \pmod{11} \\&\equiv 64 \pmod{11} \\&\equiv 9 \pmod{11}\end{aligned}$$

So Bob sends $B = 9$ to Alice.

(d) Alice computes

$$\begin{aligned} K &\equiv B^a \pmod{11} \\ &\equiv 9^8 \pmod{11} \\ &\equiv 3 \pmod{11}, \end{aligned}$$

while Bob computes

$$\begin{aligned} K &\equiv A^b \pmod{11} \\ &\equiv 3^6 \pmod{11} \\ &\equiv 3 \pmod{11}. \end{aligned}$$

This verifies that their shared key is $K = 3$.

2. (a) Bob enciphers m to

$$\begin{aligned} c &\equiv mK \pmod{p} \\ &\equiv 5 \times 3 \pmod{11} \\ &\equiv 4 \pmod{11}. \end{aligned}$$

So Bob sends $c = 4$.

(b) Alice has to calculate $cK^{-1} \pmod{p}$, where K^{-1} is the inverse of $K \pmod{11}$.

It is clear that with $K = 3$, we have $K^{-1} = 4$ since $3 \times 4 \equiv 1 \pmod{11}$. So Alice calculates

$$\begin{aligned} cK^{-1} \pmod{p} &\equiv 4 \times 4 \pmod{11} \\ &\equiv 5 \pmod{11} \end{aligned}$$

and recovers $m = 5$.

3. (a) $n = 19 \times 13 = 247$

(b) Now $\phi(n) = (19 - 1)(13 - 1) = 216$, and so Alice's private key d is the inverse of $e = 5$ in \mathbb{Z}_{216} . Using Euclid's Algorithm, we get $d = 173$.

(c) Alice calculates

$$s \equiv 93^{173} \pmod{247}.$$

Using fast exponentiation $s = 175$.

4. The easiest way to solve these problems is to run through all possible solutions.

Recall that $x - a$ is a factor of $f(x)$ if and only if $f(a) = 0$ (The Factor Theorem). So run through all values of a in $\{0, 1, 2, \dots, n - 1\}$ to see whether $f(a) = 0$. In this way, we obtain the following:

(a) The roots of $x^2 + 3x + 2$ in $\mathbb{Z}_5[x]$ are 3 and 4. Thus

$$x^2 + 3x + 2 = (x - 3)(x - 4) = (x + 2)(x + 1).$$

(b) The roots of $x^2 + 3x + 2$ in $\mathbb{Z}_7[x]$ are 5 and 6. Thus

$$x^2 + 3x + 2 = (x - 5)(x - 6) = (x + 2)(x + 1).$$

(c) The roots of $x^4 + 4$ in $\mathbb{Z}_5[x]$ are 1, 2, 3, and 4. Thus

$$x^4 + 4 = (x - 1)(x - 2)(x - 3)(x - 4) = (x + 4)(x + 3)(x + 2)(x + 1).$$

5. Putting successively $x = 0, 1, 2, \dots, 11$, we find that if $x \in \{2, 7, 10, 11\}$, then $f(x) = 0$ but, if $x \in \{0, 1, 3, 4, 5, 6, 8, 9\}$, then $f(x) \neq 0$. So the roots of $f(x)$ in $\mathbb{Z}_{12}[x]$ are 2, 7, 10, 11. This gives two distinct factorisations:

$$f(x) = (x - 2)(x - 7) = (x + 10)(x + 3)$$

and

$$f(x) = (x - 10)(x - 11) = (x + 2)(x + 1).$$

By Corollary 6.7, if p is prime, then a polynomial of degree two in $\mathbb{Z}_p[x]$ has at most two roots, and thus a unique factorisation. But if m is composite, then a polynomial of degree two in $\mathbb{Z}_m[x]$ may have more than two roots and thus more than one factorisation.