

**Lab Quiz 4**

Quiz relates to Lectures 11 and 12. Questions might have been seen in a different order on LEARN.

**QUESTION 1**

The Merkle-Damgård construction for hash functions makes use of a compression function,  $h$ , which acts on successive message blocks. A benefit of this construction is:

- (a) computation of a hash value requires a fixed number of calls to  $h$ , independent of the length of the input message
- (b) if  $h$  is collision-resistant then the whole hash function is collision-resistant
- (c) no padding is required for the input message, no matter what is the output size of  $h$
- (d) the length of the input message does not need to be included

---

---

if  $h$  is collision-resistant then the whole hash function is collision-resistant

---

---

**QUESTION 2**

Due to the birthday paradox, we can expect to find a collision in the SHA-256 hash function after around:

- (a)  $2^7$  trials
- (b)  $2^8$  trials
- (c)  $2^{128}$  trials
- (d)  $2^{255}$  trials

---

---

$2^{128}$  trials

---

---

### QUESTION 3

Suppose that an attacker has the ability to compute the output of a certain hash function for  $2^{128}$  input values. In order to prevent the attacker from finding a collision in the hash function, the output of the hash function should be of length at least:

- (a) 128 bits
- (b) 256 bits
- (c) 384 bits
- (d) 512 bits

---

---

384 bits

---

---

### QUESTION 4

A message authentication code (MAC) takes as input a message and a key and outputs a tag. To be considered secure a MAC should have the property:

- (a) the correct tag for a new message cannot be computed without the key
- (b) the message used to compute the tag cannot be distinguished from a random message
- (c) different tags are computed if a message is repeated
- (d) any output tag cannot be distinguished from a random string

---

---

the correct tag for a new message cannot be computed without the key

---

---

### QUESTION 5

Which of the following block cipher modes of operation is not designed to provide data confidentiality?

- (a) Counter mode (CTR)
- (b) Cipher block chaining (CBC)
- (c) Cipher-based MAC (CMAC)
- (d) Counter with CBC-MAC (CCM)

---

---

Cipher-based MAC (CMAC)

---

---

### **QUESTION 6**

Which of the following block cipher modes of operation is not designed to provide data integrity?

- (a) Galois counter mode (GCM)
- (b) Cipher block chaining (CBC)
- (c) Cipher-based MAC (CMAC)
- (d) Counter with CBC-MAC (CCM)

---

---

Cipher block chaining (CBC)

---

---

### **QUESTION 7**

When public key cryptography is used for encryption:

- (a) the public key of the sender is required in order to decrypt the ciphertext
- (b) the public key of the receiver is required in order to decrypt the ciphertext
- (c) the private key of the sender is required in order to decrypt the ciphertext
- (d) the private key of the receiver is required in order to decrypt the ciphertext

---

---

the private key of the receiver is required in order to decrypt the ciphertext

---

---

### **QUESTION 8**

The keys for the RSA encryption algorithm include a public exponent  $e$ , a private exponent  $d$ , and a public modulus  $n$ . It is common to choose:

- (a)  $d = 2^{16} + 1$
- (b)  $e = 2^{16} + 1$
- (c)  $e = n - 1$
- (d)  $d = n - 1$

---

---

$e = 2^{16} + 1$

---

---

**QUESTION 9**

For the RSA encryption scheme a large modulus  $n$  is chosen, typically around 2048 bits in practice. To improve efficiency, this is often used together with:

- (a) a small value for  $e$
- (b) a small value for  $d$
- (c) a small value for one of the factors of  $n$
- (d) a small value for the Euler function  $\phi(n)$

---

---

a small value for  $e$

---

---

**QUESTION 10**

For any given values  $x$  and  $m$ , the square-and-multiply algorithm when used to compute  $x^{66} \bmod m$  requires:

- (a) 5 squarings and 3 multiplications modulo  $m$
- (b) 6 squarings and 1 multiplication modulo  $m$
- (c) 8 squarings and 1 multiplication modulo  $m$
- (d) 63 squarings and 3 multiplication modulo  $m$

---

---

6 squarings and 1 multiplication modulo  $m$

---

---