SYSTEM PROGRAMMING

WEEK 5: SYSTEM DATA FILES AND PROCESS ENVIRONMENT

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Password, Group, and other system files

introduction

A Unix system requires numerous data files for normal operations

- /etc/passwd for user log-in or ls -l, etc
- /etc/group

These data files have been ASCII text files and were read with the standard I/O libaray.

 sequential scan of such files became time consuming for a large system

We want to other data format other than ASCII, but still provide an interface for any application.

PASSWORD, GROUP, AND OTHER SYSTEM FILES: PASSWORD FILES

Password File

The UNIX System's password file <pwd.h>, called the user database by POSIX.1, contains the following fields

Description	struct passwd member			POSIX.1	FreeBSD 8.0	Linux 3.2.0	Mac OS X 10.6.8	Solaris 10
user name	char	*pw	name	•	•	•	•	•
encrypted password	char	*pw	passwd		•	•	•	•
numerical user ID	uid_t	pw	uid	•	•	•	•	•
numerical group ID	gid_t	pw	gid	•	•	•	•	•
comment field	char	*pw	gecos		•	•	•	•
initial working directory	char	*pw	dir	•	•	•	•	•
initial shell (user program)	char	*pw	shell	•	•	•	•	•
user access class	char	*pw	class		•		•	
next time to change password	time_t	pw_	change		•		•	
account expiration time			expire		•		•	

Figure: Fileds in /etc/passwd file

On your machine, you can see the contents by typing cat /etc/passwd

- 1 root:x:0:0:root:/root:/bin/bash
- squid:x:23:23::/var/spool/squid:/dev/null
- 3 nobody:x:65534:65534:Nobody:/home:/bin/sh
- sar:x:205:105:Stephen Rago:/home/sar:/bin/bash



Password File cnt'd

- if ecrypted passwd field is empty, it usually means that the user does not have passwd
- shell field contains the name of a program to be used as the login shell for the user. If this field is empty, the default is /bin/sh
- /dev/null, /bin/false, or /bin/true in the shell field prevents a particular user from logging in to a system
- nobody user name can be used to allow people to log, but with a user ID (65534) and group ID (65534)

Password File API

```
#include <pwd.h>
struct passwd *getpwuid(uid_t uid);
struct passwd *getpwnam(const char *name);
// Both return: pointer if OK, NULL on error
```

getpwuid function is used by ls(1) program to map the numberical user ID contained in an i-node into a user's login name

getpwnam function is used by the login(1) program when we enter our login name

Both functions return a pointer to passwd structure that the functions fill in

They are only for looking up either a login name or a user ID

Password File API cnt'd

```
#include <pwd.h>
struct passwd *getpwent(void);
// Returns: pointer if OK, NULL on error or end of file
void setpwent(void);
void endpwent(void);
```

getpwent to return the next entry in the password file.

The function setpwent rewinds whatever files it uses, and endpwent closes these files.

When using getpwent, we must always be sure to close these files by calling endpwent when we're through.

getpwnam implementation

```
#include <pwd.h>
    #include <stddef.h>
    #include <string.h>
    struct passwd *
    getpwnam(const char *name)
       struct passwd *ptr:
      setpwent():
10
      while ((ptr = getpwent()) != NULL)
11
         if (strcmp(name, ptr->pw_name) == 0)
12
            break: /* found a match */
       endpwent():
14
       return(ptr); /* ptr is NULL if no match found */
15
16
```

more interested reader can go on and have a look at getspnam and setspent for shadow password files



PASSWORD, GROUP, AND OTHER SYSTEM FILES: GROUP FILES

Group File

The UNIX System's group file, called the group database by POSIX.1, contains the fields shown in the figure. These fields are contained in a group structure that is defined in <grp.h>.

Description	struct group member		POSIX.1	FreeBSD 8.0	Linux 3.2.0	Mac OS X 10.6.8	Solaris 10
group name encrypted password numerical group ID array of pointers to individual user names	char char int char	gr_gid	:	:	:	•	:

Figure: Fields in /etc/group file

The field gr_mem is an array of pointers to the user names that belong to this group. This array is terminated by a null pointer.

Group File API

We can look up either a group name or a numerical group ID with the following

```
#include <grp.h>
struct group *getgrgid(gid_t gid);
struct group *getgrnam(const char *name);
// Both return: pointer if OK, NULL on error
```

Like the password file functions, both of these functions normally return pointers to a static variable, which is overwritten on each call.

Followings are used for searching the entire group file



PASSWORD, GROUP, AND OTHER SYSTEM FILES: OTHERS SYSTEM

DATABASES FILES

Other System Database Files

There are other system files and has similar routines for accessing them

Description	Data file	Header	Structure	Additional keyed lookup functions
passwords groups shadow	/etc/passwd /etc/group /etc/shadow	<pre><pwd.h> <grp.h> <shadow.h></shadow.h></grp.h></pwd.h></pre>	passwd group spwd	getpwnam, getpwuid getgrnam, getgrgid getspnam
hosts networks protocols services	/etc/hosts /etc/networks /etc/protocols /etc/services	<netdb.h> <netdb.h> <netdb.h> <netdb.h></netdb.h></netdb.h></netdb.h></netdb.h>	hostent netent protoent servent	getnameinfo, getaddrinfo getnetbyname, getnetbyaddr getprotobyname, getprotobynumber getservbyname, getservbyport

Figure: Similar routines for accessing system data files

- **get** function reads the enxt record, opening the file if necessary
- set function that opens the file, if not already open, and rewinds the file
- end entry that closes the data file. We always have to call this function when we're done.



System Identification

POSIX.1 defines the uname function to return information on the current host and operating system

```
#include <sys/utsname.h>
int uname(struct utsname *name);
// Returns: non-negative value if OK, 1 on error
```

- We pass the address of a utsname structure to this function, and the function then fills it in.
- The structure contains fields like sysname, nodename, release, version, and machine; but, they depend on the implementation
- O This is used by the shell command uname(1)

```
#include <unistd.h>
int gethostname(char *name, int namelen);
// Returns: 0 if OK, 1 on error
```

Historically, BSD-derived systems provided the gethostname function to return only the name of the host

SIL



Password, Group, and other <u>system</u> files: Time and Date

Time and Date Routines

The basic time service provided by the UNIX kernel counts the number of seconds that have passed since the Epoch: 00:00:00 January 1, 1970, Coordinated Universal Time (UTC).

- they are represented in a time_t data type
- they are called *calendar times* and represent both the time and the date

The Unix System has always differed from other operating systems in

- keeping time in UTC instead of the local time
- automatically handling conversions, such as daylight saving time,
- keeping the time and date as a single quantity.

Time and Date Routines API

The time function returns the current time and date

```
#include <time.h>
time_t time(time_t *calptr);
// Returns: value of time if OK, 1 on error
```

- The time value is always returned as the value of the function.
- If the argument is non-null, the time value is also stored at the location pointed to by *calptr*

Time and Date Routines cnt'd

The clock_gettime function can be used to get the time of the specified clock

```
#include <sys/time.h>
int clock_gettime(clockid_t clock_id, struct timespec *tsp); // get the time
int clock_getres(clockid_t clock_id, struct timespec *tsp); // resolution of a
    system clock
int clock_settime(clockid_t clock_id, const struct timespec *tsp); // set the
    time
// Returns: 0 if OK, 1 on error
```

Identifier	Option	Description
CLOCK_REALTIME		real system time
CLOCK_MONOTONIC	_POSIX_MONOTONIC_CLOCK	real system time with no negative jumps
CLOCK_PROCESS_CPUTIME_ID	_POSIX_CPUTIME	CPU time for calling process
CLOCK_THREAD_CPUTIME_ID	_POSIX_THREAD_CPUTIME	CPU time for calling thread

Figure: Clock type identifiers



Time and Date Routines cnt'd

This function is now obsolete. However, a lot of programs stil use it, because it provides greater resolution (up to a microsecond) than the time function

```
int gettimeofday(struct timeval *restrict tp, void *restrict tzp);
// Returns: 0 always
```

Time and Date Routines cnt'd: formating

The strftime and strftime_l functions are the same, except that the strftime_l function allows the caller to specify the locale as an argument.

Relationship of Time Functions

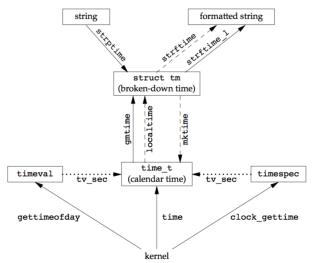


Figure: Relationship of the various time functions



Time and Date Routines cnt'd: tm structure

```
struct tm { /* a broken-down time */
   int tm_sec; /* seconds after the minute: [0 - 60] */
   int tm_min; /* minutes after the hour: [0 - 59] */
   int tm_hour; /* hours after midnight: [0 - 23] */
   int tm_mday; /* day of the month: [1 - 31] */
   int tm_mon; /* months since January: [0 - 11] */
   int tm_year; /* years since 1900 */
   int tm_yday; /* days since Sunday: [0 - 6] */
   int tm_yday; /* days since January 1: [0 - 365] */
   int tm_isdst; /* daylight saving time flag: <0, 0, >0 */
};
```

strftime example

make strftime; ./strftime

```
time t t:
       struct tm *tmp;
       char buf1[16];
       char buf2[64];
       time(&t);
       tmp = localtime(&t);
       if (strftime(buf1, 16, "time and date: %r, %a %b %d, %Y", tmp) == 0)
          printf("buffer length 16 is too small\n");
       else
10
          printf("%s\n", buf1);
11
12
       if (strftime(buf2, 64, "time and date: %r, %a %b %d, %Y", tmp) == 0)
          printf("buffer length 64 is too small\n");
13
       else
14
          printf("%s\n". buf2):
15
16
17
       exit(0);
18
```

PASSWORD, GROUP, AND OTHER SYSTEM FILES: LAST WORDS

Homework

- Download Mobibench code from https://github.com/ESOS-Lab/Mobibench/tree/master/shell
- Measured time is stored in timeval which is in microseconds.
- Change the code to measure the performance in nanoseconds.
- \bigcirc use diff to create patch file and show your work.
- print out the patch file and handed in on the first class of 2016-10-12

How to create a patch file

The original file foo.c

```
#include <stdio.h>
int main() {
printf("Hello World\n");
}
```

Your modified file foobar.c

```
#include <stdio.h>

int main(int argc, char *argv[]) {
 printf("Hello World\n");
 return 0;
}
```

Create the patch file using diff command diff -u foo.c foobar.c > foobar.patch

```
1  --- foo.c 2016-09-30 17:04:42.000000000 +0900
2  +++ foobar.c 2016-09-30 17:05:03.000000000 +0900
3  @@ -1,5 +1,6 @@
4  -#include <stdio.h>
5  +#include <stdio.h>
6
7  -int main() {
8  +int main(int argc, char *argv[]) {
9    printf("Hello World\n");
10  -}
11  +return 0;
12  +}
```



5

PROCESS ENVIRONMENT

Introduction

In this chapter, we'll see

- how the main function is called when the program is executed
- called when the program is executed
- how command-line arguments are passed to the new program
- what the typical memory layout looks like
- how to allocate additional memory
- how the process can use environment variables
- $\, \bigcirc \,$ and various ways for the process to terminate
- the longjmp and setjmp functions and their interaction with the stack

We finish the chapter by examining the resource limits of a process.

PROCESS ENVIRONMENT: PROCESS

CREATION AND TERMINATION

main Function

```
int main(int argc, char *argv[]);
```

- *argc*: the number of commands-line arguments
- *argv*: an array of pointers to the arguments

The executable program file specifies this routine as the starting address for the program;

- this is set up by the link editor when it is invoked by the C compiler.
- This start-up routine takes values from the kernel—the command-line arguments and the environment—and sets things up

Process Termination

There are eight ways for a process to terminate

Normal termination

- Return from main
- Calling exit
- Calling _exit or _Exit
- O Return of the last thread from its start routine
- \odot Calling <code>pthread_exit</code> from the last thread

Abnormal termination (We'll discuss them later)

- calling abort
- Receipt of a signal
- Response of the last thread to a cancellation request

exit(3) Functions

Following functions terminate a program normally

```
#include <stdlib.h>
void exit(int status); // return to the kernel after some cleanup
void _Exit(int status); // return to the kernel immediately

#include <unistd.h>
void _exit(int status); // return to the kernel immediately
```

exit(3) Functions

All three exit functions expect a single integer argument, which we call the exit status. E.g., exit(0);

Most UNIX System shells provide a way to examine the exit status of a process.

The exit status of the process is undefined, if

- any of these functions is called without an exit status,
- main does a return without a return value, or
- the main function is not declared to return an integer,

exit(3) function example

The classic "hello, world" example

```
#include <stdio.h>
    #include <stdlib.h>
    #include <unistd.h>
    int mv main()
       printf("hello, world\n");
       sleep(5);
       return 0;
11
    void _start()
12
13
       int ret = mv main():
14
       exit(ret);
15
16
```

```
$ make hello
$ ./hello
hello, world
^C
$ echo $?
130
$? prints exit status of previous
process
more on exit status http://tldp.org/
LDP/abs/html/exitcodes.html
```

atexit(3) function

With ISO C, a process can register at least 32 functions that are automatically called by exit. These are called *exit handlers* and are registered by calling the atexit function.

```
#include <stdlib.h>
int atexit(void (*func)(void));
// Returns: 0 if OK, nonzero on error
```

How Process are Created and Terminated

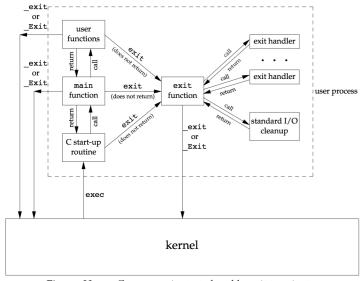


Figure: How a C program is started and how it terminates



PROCESS ENVIRONMENT: ENVI-

RONMENTS

Command-Line Arguments

When a program is executed, the process that does the exec can pass command-line arguments to the new program.

cd codes; make echoall

```
#include <stdio.h>
#include <stdio.h>
#include <stdlib.h>
#int
main(int argc, char *argv[])

{
    int i;
    for (i = 0; i < argc; i++)
        /* echo all command-line args */
        printf("argv[%d]: %s\n", i, argv[i]);
    exit(0);
}</pre>
```

try, ./echoall arguments tests

This you might want to test

- O How many arguments can you pass
- how can you insert multiple lines of text as arguments



Environment List

Each program is also passed an environment list

- it is an array of character pointers
- each pointer containing the address of a null-terminated C string
- The address of the array of pointers is contained in the global variable environ extern char **environ;

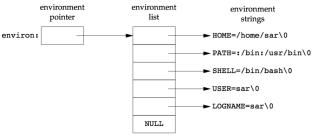


Figure: Environment consisting of five C character strings

Environment Variables

Variable	POSIX.1	FreeBSD Linux Mac OS X Solaris				Description	
		8.0	3.2.0	10.6.8	10	Description	
COLUMNS	•	•	•	•	•	terminal width	
DATEMSK	XSI		•	•	•	getdate(3) template file pathname	
HOME	•	•	•	•	•	home directory	
LANG	•	•	•	•	•	name of locale	
LC_ALL	•	•	•	•	•	name of locale	
LC_COLLATE	•	•	•	•	•	name of locale for collation	
LC_CTYPE	•	•	•	•	•	name of locale for character classification	
LC_MESSAGES	•	•	•	•	•	name of locale for messages	
LC_MONETARY	•	•	•	•	•	name of locale for monetary editing	
LC_NUMERIC	•	•	•	•	•	name of locale for numeric editing	
LC_TIME	•	•	•	•	•	name of locale for date/time formatting	
LINES	•	•	•	•	•	terminal height	
LOGNAME	•	•	•	•	•	login name	
MSGVERB	XSI	•	•	•	•	fmtmsg(3) message components to process	
NLSPATH	•	•	•	•	•	sequence of templates for message catalogs	
PATH	•	•	•	•	•	list of path prefixes to search for executable file	
PWD	•	•	•	•	•	absolute pathname of current working directory	
SHELL	•	•	•	•	•	name of user's preferred shell	
TERM	•	•	•	•	•	terminal type	
TMPDIR	•	•	•	•	•	pathname of directory for creating temporary files	
TZ	•	•	•	•	•	time zone information	

Figure: Environment consisting of five C character strings



Environment List cnt'd

ISO C specifies that the main function be written with two arguments POSIX.1 specifies that environ should be used instead of the (possible) third argument

- Access to specific environment variables is normally through the getenv and putenv functions
- But to go through the entire environment, the environ pointer must be used

Environment List APIs

```
#include <stdlib.h>
char *getenv(const char *name);
// Returns: pointer to value associated with name, NULL if not found
int putenv(char *str);
// Returns: 0 if OK, nonzero on error
int setenv(const char *name, const char *value, int rewrite);
int unsetenv(const char *name);
// Both return: 0 if OK, 1 on error
```

Fuction	ISO C	POSIX.1	FreeBSD 8.0	Linux 3.2.0	Mac OS X 10.6.8	Solaris 10
getenv	•	•	•	•	•	•
putenv		XSI	•	•	•	•
setenv		•	•	•	•	
unsetenv		•	•	•	•	
clearenv				•		

Table: Support for various environment list functions

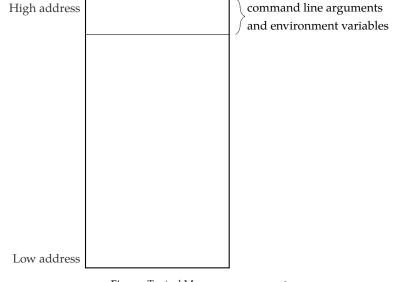
PROCESS ENVIRONMENT : MEMORY

The Pieces of a C program

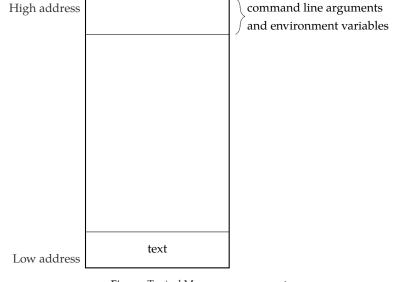
- **Text segment** Machine instructions for CPU. It is sharable. It is **read-only** to prevent from accidental modification
- Initialized data segment Or, data segment. Contains variables that
 are specifically initialized in the program
 int maxcount = 99;
- Uninitialized data segment Often called the "bss" segment (A.K.A
 "block started by symbol")
 long sum[1000];
- **Stack** Automatic variables are stored. Each time a function is called, the address of where to return to and info about the caller's environment (i.e., registers) are saved on the stack.
- **Heap** dynamic memory allocation.

High address Low address

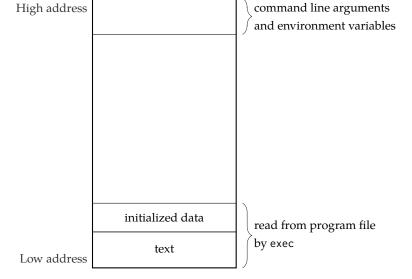




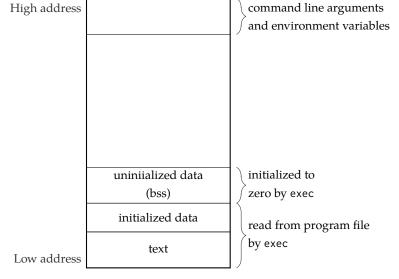




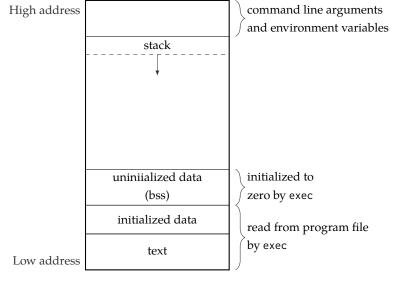




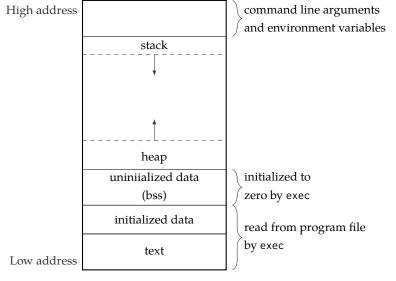














Shared Libraries

size(1) command reports the size (in bytes) of the text, data, and bss segments.

Mac does not support static linking of a library, but on Linux supports static linking.

Compare the size of a program before and after static linking.

Dynamic linking

```
cc -o layout layout.c
file layout
ldd layout
size layout
objdump -d layout > obj_layout
wc -l obj_layout
```

Static linking

```
cc -static -o layout_st layout.c
file layout_st
ldd layout_st
size layout_st
objdump -d layout_st > obj_layout_st
wc -l obj_layout_st
```

Memory Allocation

```
#include <stdlib.h>
void *malloc(size_t size);
void *calloc(size_t nobj, size_t size);
void *realloc(void *ptr, size_t newsize);
// All three return: non-null pointer if OK, NULL on error
void free(void *ptr);
```

- **malloc** allocates a specified number of bytes of memory. The initial value of the memory is indeterminate
- **calloc** allocates a specified number of bytes of memory and initializes them to o bits
- realloc increases or decreases the size of a preivously allocated area
 free deallocates the area. Do not free what is already freed

PROCESS ENVIRONMENT: MISC.

setjmp and longjmp Functions

In C, we can't goto a label that's in another function. Instead, we must use the setjmp and longjmp functions to perform this type of branching

```
#include <setjmp.h>
int setjmp(jmp_buf env);
// Returns: 0 if called directly, nonzero if returning from a call to longjmp
void longjmp(jmp_buf env, int val);
```

- call setjmp from the location that we want to return to.
 - jmp_buf data type is some form of array that is capable of holding all the information required to restore the status of the stack to the state when we call longjmp
- when we encounter an error, we call longjmp with two arguments
 - ∘ jmp_buf env used in setjmp
 - $\circ\,$ val is nonzero value that becomes the return value from setjmp

Last Words

○ read chapter 8