# SYSTEM PROGRAMMING

#### **WEEK 7: PROCESS RELATIONSHIPS**

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Updated: 2016-10-19

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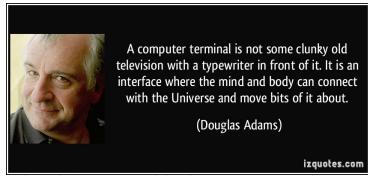


#### introduction

This chapter covers following items

- Terminal Logins
- Sessions
- Controlling Terminal
- Job Control

#### **Terminal**



**Figure:** Douglas Noel Adams was an English author, scriptwriter, essayist, humorist, satirist and dramatist, best known as the author of The Hitchhiker's Guide to the Galaxy" 1978

In early Unix systems

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- The terminals were either local or remote
- Users login through a terminal device driver in the kernel
- A host had a fixed number of terminal devices type who in the shell

```
James console Oct 12 15:36
James ttys000 Oct 13 22:02
James ttys001 Oct 13 22:08
```

Mac OS X and Linux login procedure follows essentially the same steps as the BSD

The system administrator creates /etc/ttys, ttys(5), that has one line per terminal device

Each line specifies the name of the device and other parameters that are passed to the getty(8) program

 $process\ ID\ {\bf 1}$ 

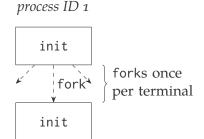
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- 3. creates empty environment

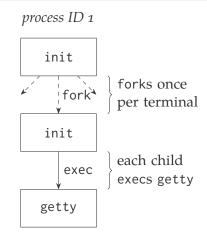
#### process ID 1

init

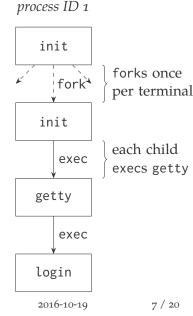
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- 7. reads user name
- 8. initial environment set



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- followed by exec of the program login





init(8) PID 1 PPID 0 EUID 0

reads/etc/ttys

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opens terminal
prints "login:"
read username

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getpass(3), encrypt, compare with getpwnam(3) register login in system database

read/display vaious files

initgroups(3)/setgid(2), initialize environment

chdir(2) to home directory

chown(2) terminal device

setuid(2) to user's uid, exec(3) shell



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init(8)	PID 1	PPID o	EUID o
reads/etc/ttys	3		
getty(8)	PID N	PPID o	EUID o
opens terminal			
prints "login:"			
read username			
login(1)	PID N	PPID o	EUID o
<pre>getpass(3), encrypt, compare with getpwnam(3)</pre>			
register login in system database			
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<pre>initgroups(3)/setgid(2), initialize environment</pre>			
chdir(2) to home directory			
chown(2) terminal device			
setuid(2) to user's uid, exec(3) shell			
\$SHELL	PID N	PPID o	EUID U
ls(1)	PID M	PPID N	EUID U
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### **Process Groups**

#### Each process belongs to a process group

- it is a collection of one or more processes
- usually associated with the same job
- the group has a unique process group ID
- $\odot$  the process group exista as long as one process is in the group

```
#include <unistd.h>
pid_t getpgrp(void);
// Returns: process group ID of calling process
```

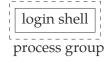


# Process group cont'd

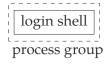
A process joins an existing process group or creates a new process group by calling setpgid

```
#include <unistd.h>
int setpgid(pid_t pid, pid_t pgid);
// Returns: 0 if OK, 1 on error
```

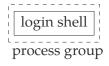
- sets the process group ID to pgid in the process whose process ID equals to pid
- if pgid == pid, then pid == process group leader
- if pid == o, caller process ID is used
- if pgid == o, group ID == pid
- O A porcess can set the process group ID of itself or its children







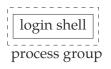
proc1

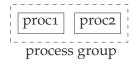


proc1 proc2

The processes in a process group are usually placed ther by a shell pipeline

proc1 | proc2 &



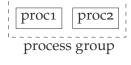


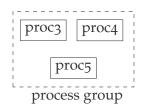


The processes in a process group are usually placed ther by a shell pipeline

```
proc1 | proc2 &
proc3 | proc4 | proc5
```









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```
proc1 | proc2 &
proc3 | proc4 | proc5
```



session

Session is a collection of one or more process groups



A process establishes a new session by calling the setsid function

```
#include <unistd.h>
pid_t setsid(void);
// Returns: process group ID if OK, 1 on error
```

If the calling process is not a process group leader, this fuction creates a new session. Three things happen

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- 2. the process becomes the process group leader (pgid ==pid)
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- 1. process becomes the *session leader*, and is only process in this new session
- 2. the process becomes the process group leader (pgid ==pid)
  - if the caller is already a process group leader, then returns an error
- 3. No contorlling terminal

getsid function returns the process group ID of a process's session leader

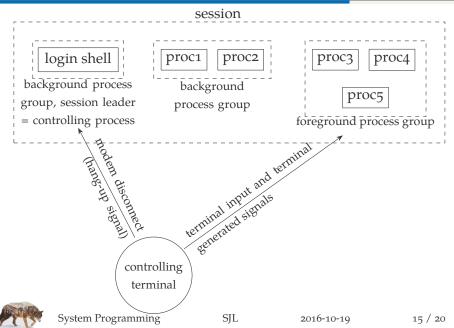
```
#include <unistd.h>
pid_t getsid(pid_t pid);
// Returns: session leader's process group ID if OK, 1 on error
```

if pid == 0, it returns the pgid of calling process's session leader

### Controlling terminal

- Session can have a single controlling terminal
- session leader that connects to controlling terminal is controlling process
- process groups are divided into a single forground process group and one or more background process groups
- o interrupt signals are sent to foreground process group

# Controlling Terminal



### Job Control

We can start a job in either the forground or the background

foreground vi main.c starts a job in the foreground background make all & start a job in the background

```
$ make all > Make.out &
[1] 1475
$ pr *.c | lpr &
[2] 1490
$ just press RETURN
[2] + Done pr *.c | lpr &
[1] + Done make all > Make.out &
```



### lob Control cont'd

The foreground jobs are affected by some special characters, which generate signals

- Interrupt character (typically DELETE or Control-C) generates SIGINT
- Quit character (typically Control-bashslash) generates SIGQUIT
- Suspend character (typically Control-Z) generatea SIGTSTP

### Job Control cont'd

```
1 $ cat > temp.foo &
    [1] 1681
2 $
    [1] + Stopped (SIGTTIN) cat > temp.foo &
3 $ fg %1
4 cat > temp.foo
5 hello, world
6 ^D
7 $ cat temp.foo
hello, world
```

- 1. start in background, but it'll read from standard input
- 2. we press RETURN
- 3. bring job number 1 into the foreground
- 4. the shell tells us which job is now in the foreground
- 5. enter one line
- 6. type the end-of-file character
- 7. check that the one line was put into the file



# LAST WORDS

### **Last Words**

O Read chapter 10

