

# Problem Solving: Combinatorics

## Problem-01: Only Pluses

(Problem Source: <https://codeforces.com/problemset/problem/1992/A> )

Kmes has written three integers  $a$ ,  $b$  and  $c$  in order to remember that he has to give Noobish\_Monk  $a \times b \times c$  bananas. Noobish\_Monk has found these integers and decided to do the following **at most 5 times**:

- pick one of these integers;
- increase it by 1.

For example, if  $a=2$ ,  $b=3$  and  $c=4$ , then one can increase  $a$  three times by one and increase  $b$  two times. After that  $a=5$ ,  $b=5$ ,  $c=4$ . Then the total number of bananas will be  $5 \times 5 \times 4 = 100$ .

What is the maximum value of  $a \times b \times c$  Noobish\_Monk can achieve with these operations?

## Problem-01 (Cont.)

### Input

Each test contains multiple test cases. The first line of input contains a single integer  $t$  ( $1 \leq t \leq 1000$ ) — the number of test cases. The description of the test cases follows.

The first and only line of each test case contains three integers  $a$ ,  $b$  and  $c$  ( $1 \leq a, b, c \leq 10$ ) — Kmes's integers.

### Output

For each test case, output a single integer — the maximum amount of bananas Noobish\_Monk can get.

### Example

#### Input

2

2 3 4

10 1 10

#### Output

100

600

# The Solution

## **Input:**

- An integer  $t \rightarrow$  number of test cases
- For each test case: three integers  $a, b, c$

## **Output:**

- For each test case, the product  $a * b * c$  after incrementing the smallest number 5 times

## **Steps**

- 1. Read** the number of test cases  $t$ .
- 2. Repeat** the following for each test case:
  - 1. Read** three integers  $a, b, c$ .
  - 2. Repeat 5 times:**
    - If  $a$  is less than or equal to both  $b$  and  $c$ , increment  $a$  by 1.
    - Else if  $b$  is less than or equal to both  $a$  and  $c$ , increment  $b$  by 1.
    - Else increment  $c$  by 1.
  - 3. Compute** the product  $a * b * c$ .
  - 4. Print** the result.

```
#include<iostream>
using namespace std;
int main()
{
    int t;
    cin>>t;
    while(t--)
    {
        int a,b,c;
        cin>>a>>b>>c;
        for(int i=0;i<5;i++)
        {
            if(a<=b && a<=c)
                a++;
            else if(b<=c && b<=a)
                b++;
            else
                c++;
        }
        cout<<a*b*c;
    }
}
```

What will be its  
Complexity?

## Problem-02: trailing zeroes in the factorial of n

Given an integer  $n$ , determine the number of trailing zeroes in the factorial of  $n$  ( $n!$ ).

A trailing zero is defined as a zero that appears at the end of a number, after the last non-zero digit. For example, 1200 has two trailing zeroes.

You must design an efficient algorithm that avoids direct computation of  $n!$ , since factorial values grow extremely large for even moderate values of  $n$ . Instead, focus on analyzing the factors that contribute to trailing zeroes.

## Problem-02 (Cont.)

### Input

A single integer  $n$  ( $1 \leq n \leq 10^9$ )

### Output

An integer representing the number of trailing zeroes in  $n!$

### Example

Input	Output
10	2

Explanation:  $10! = 3,628,800$ , which has **2 trailing zeroes**.



# The Solution

A factorial, denoted as  $n!$ , is the product of all positive integers from 1 to  $n$ . Factorials grow very quickly, and for larger values of  $n$ , the result often ends with several trailing zeroes (zeros at the end of the number).

For example:

$5! = 120 \rightarrow$  has 1 trailing zero

$10! = 3,628,800 \rightarrow$  has 2 trailing zeroes

### **Why do trailing zeroes appear?**

- A trailing zero is produced when a number is divisible by 10.
- $10 = 2 \times 5$ .
- In factorials, there are always more factors of 2 than 5.
- So, the number of trailing zeroes depends only on the number of times 5 divides the numbers from 1 to  $n$ .

### **Formula:**

Trailing Zeroes in  $n! = \lfloor n/5 \rfloor + \lfloor n/25 \rfloor + \lfloor n/125 \rfloor + \dots$

We keep dividing  $n$  by 5 until it becomes 0.

## **Algorithm: Trailing Zeroes in Factorial**

**Input:** An integer  $n$

**Output:** Number of trailing zeroes in  $n!$

**1.Initialize** count = 0.

**2.Repeat until  $n$  becomes 0:**

- Divide  $n$  by 5 (integer division).
- Add the quotient to count.
- Update  $n = n / 5$ .

**3.Return** count.

```
#include <iostream>
using namespace std;
```

Complexity?

```
int main() {
    int n;
    cout << "Enter a number: ";
    cin >> n;

    int count = 0;

    while (n > 0) {

        n = n / 5;
        count = count+n;
    }

    cout << "Trailing zeroes in factorial = " << count << endl;

    return 0;
}
```

## Problem-03: count digits in a factorial

Given an integer  $n$ , find the number of digits that appear in its factorial, where factorial is defined as,  $\text{factorial}(n) = 1*2*3*4*.....*n$  and  $\text{factorial}(0) = 1$

*Example:*

*Input: 5*

*Output: 3*

*Explanation:  $5! = 120$ , that has, 3 digits*

*Input: 10*

*Output: 7*

*Explanation:  $10! = 3628800$ , that has, 7 digits*

# The Solution

Using logarithmic property:

We know,  $\log(a*b) = \log(a) + \log(b)$

Therefore:

$$\log(n!) = \log(1*2*3*... * n) = \log(1) + \log(2) + ..... + \log(n)$$

Now, observe that the floor value of log base 10 increased by 1, of any number, gives the number of digits present in that number.

Hence, output would be :  $\text{floor}(\log(n!)) + 1$ .

## **Algorithm: Counting Digits in a Factorial**

**Input:** An integer  $n$

**Output:** Number of decimal digits in  $n!$

**1. Handle base cases:**

If  $n = 0$  or  $n = 1$ , return 1 (since  $0! = 1! = 1$ ).

**2. Initialize sum:**

Let  $\text{sum} = 0.0$  (a floating-point variable to store the sum of logarithms).

**3. Compute log10 of factorial:**

1. For each integer  $i$  from 2 to  $n$ :
  1. Add  $\log_{10}(i)$  to  $\text{sum}$ .

**4. Calculate number of digits:**  $\text{digits} = \text{floor}(\text{sum}) + 1$

**5. Return digits**



```
#include <iostream>
#include <cmath>
using namespace std;
int main() {
    int n;
    cin >> n;

    if (n == 0 || n == 1)
    {
        cout << 1 << endl;
        return 0;
    }

    double sum = 0;
    for (int i = 2; i <= n; i++)
    {
        sum = sum + log10(i);
    }
    int digits = floor(sum) + 1;
    cout << digits << endl;
    return 0;
}
```

Complexity?