

CSE 2103

(Data Structures)

Lecture on

Chapter-6: Stacks, Queues, Recursion

By

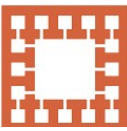
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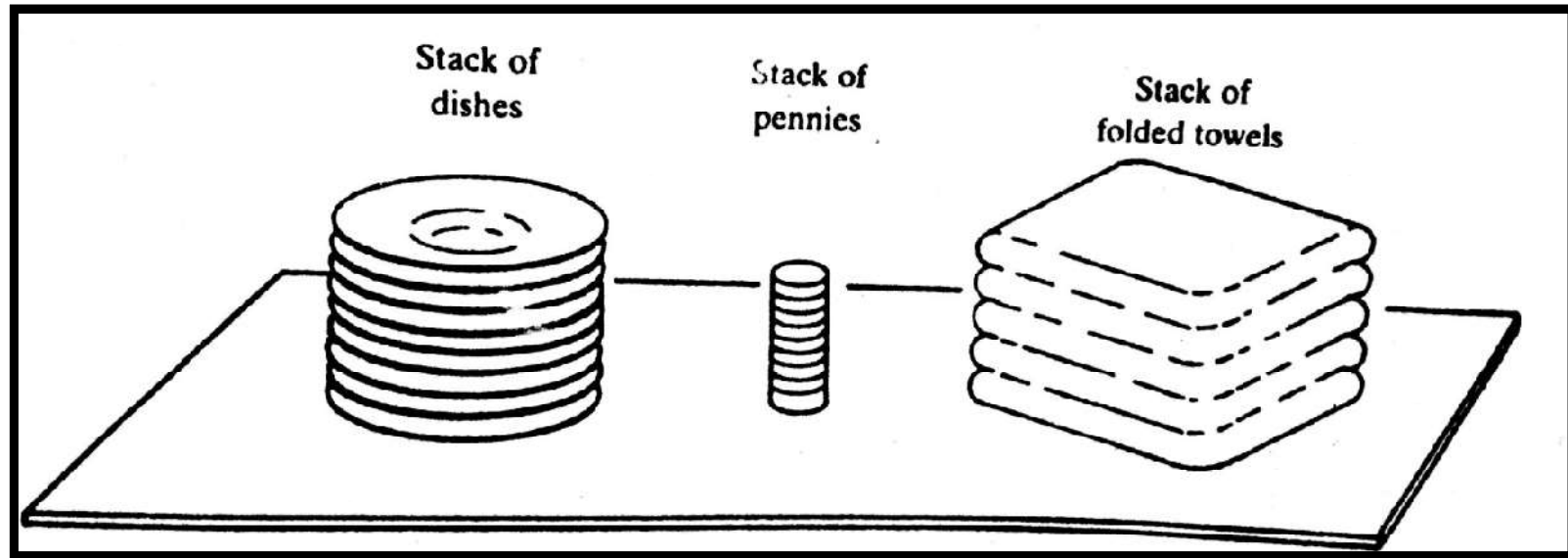


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STACKS: Introduction

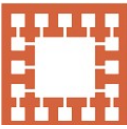


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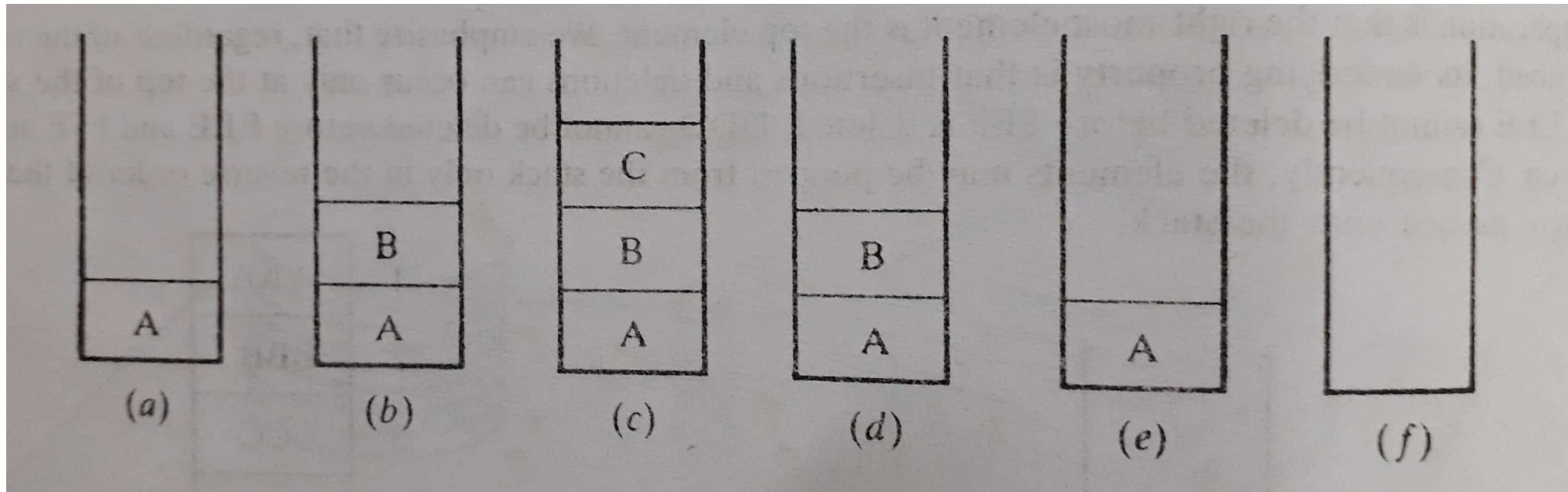
- A stack is a linear structure in which items may be added or removed only at one ends
- Stacks are also called Last-in-First-out (LIFO) list.
- Other names:
 - ❖ Piles
 - ❖ Push-down lists.
- Although the stack may be a very restricted type of data structure, it has many important applications in computer science.

STACKS: Dealing Postponed Decisions

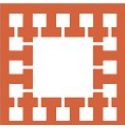


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Stacks are frequently used to indicate the order of the processing of data when certain steps of the processing must be postponed until other conditions are fulfilled.



STACKS: Special Terminology



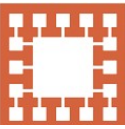
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Special terminology is used for two basic operations associated with stacks:



- (a) “Push”-used to insert an element into a stack
- (b) “Pop”-used to delete an element from a stack.

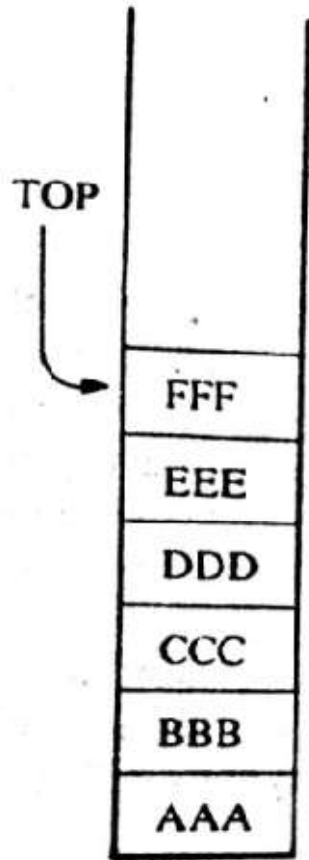
STACKS: Push Example



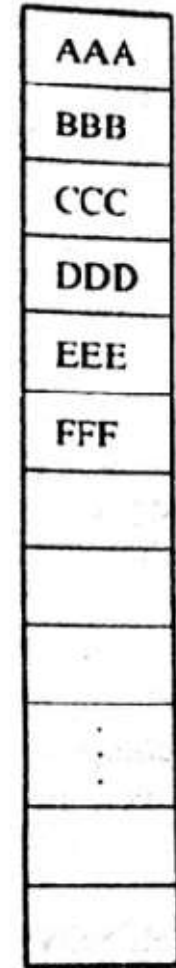
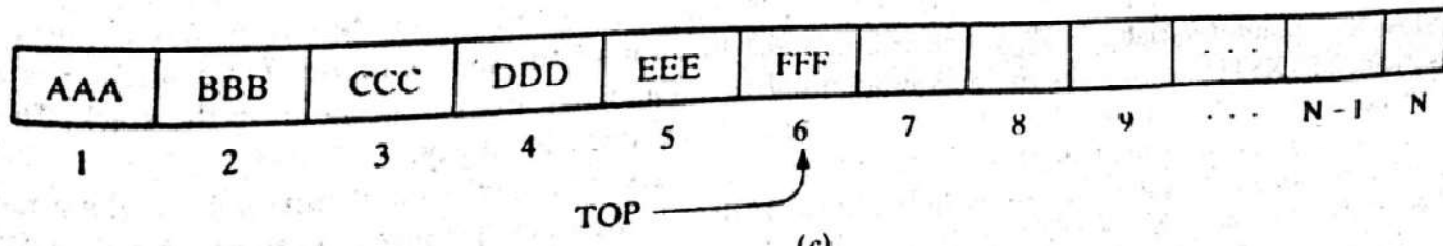
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Suppose the following 6 elements are pushed, in order, onto an empty stack:

AAA, BBB, CCC, DDD, EEE, FFF



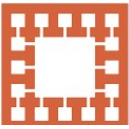
(a)



(b)

Diagram of Stack

STACKS: Array Representation



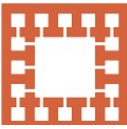
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- Stack may be represented in the computer in various ways,
 - ✓ usually by means of a one-way list, or
 - ✓ A linear array
- Usually Stacks will be maintained by.....
 - A linear array, **STACK**, (main lists)
 - A pointer variable, **TOP**, (contains the locations of the top element of the stack)
 - A variable **MAXSTK**, (gives the maximum number of elements that can be held by the stack)

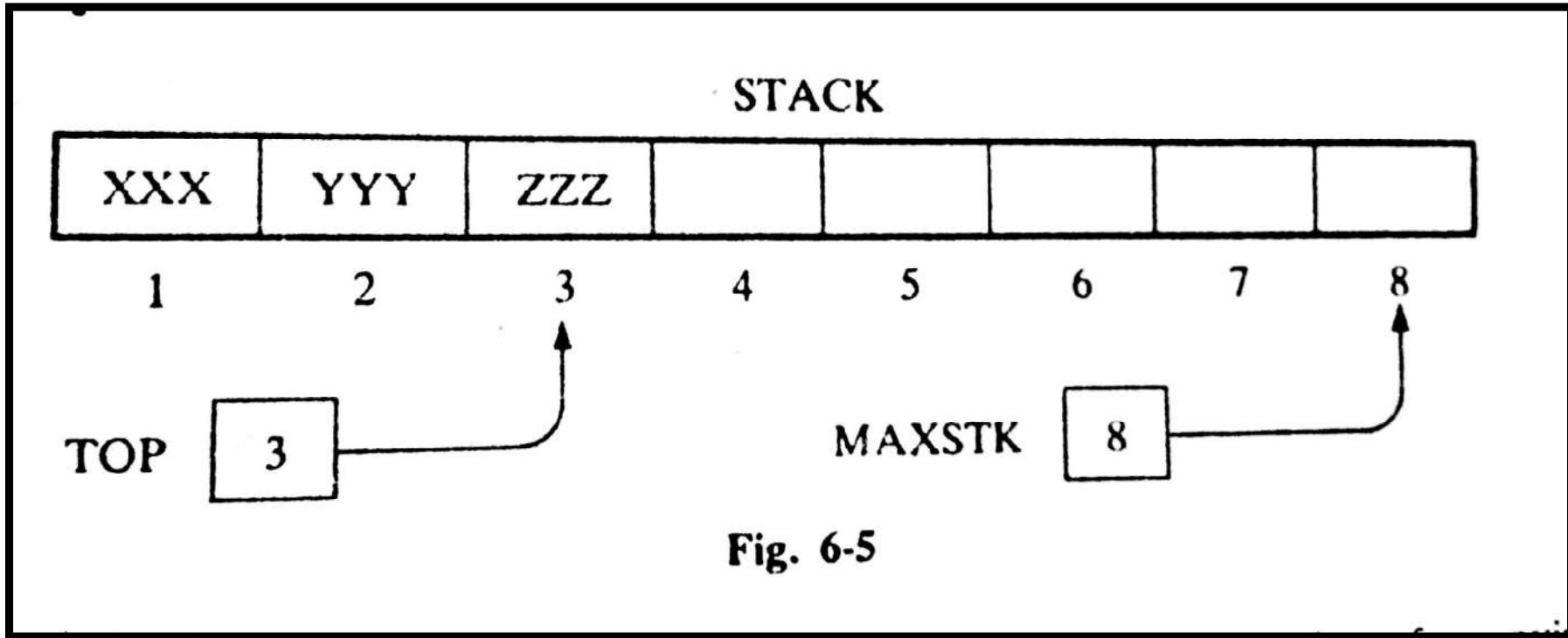
The condition **TOP=0** or **TOP=NULL** of a **STACK** indicates.....?

EMPTY

STACKS: Array Representation Example



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Present Status: TOP=3;

MAXSTK=8;

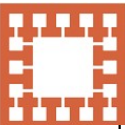
There is room for 5 more items in the stack

Future Implementation:

Pushing an item onto a stack

Popping an item from a stack

STACKS: Pushing Algorithm



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PUSH (**STACK**, **TOP**, **MAXSTK**, **ITEM**)

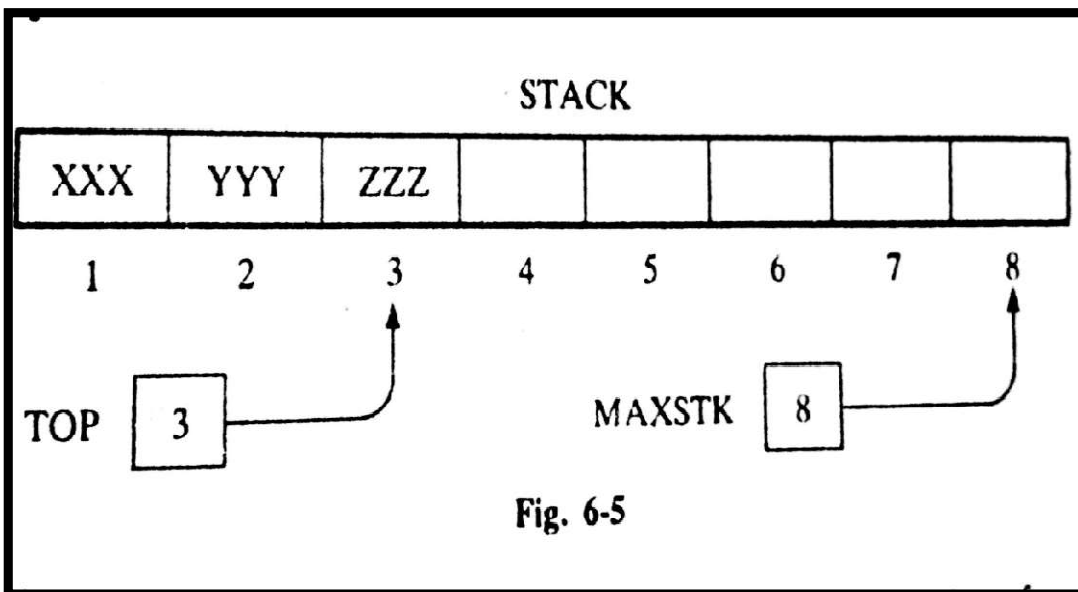
Step 1. [Stack already filled?]

If **TOP**=**MAXSTK**, then print: **OVERFLOW**, and Return.

Step 2. Set **TOP**=**TOP**+1. [Increase TOP by 1.]

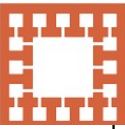
Step 3. Set **STACK**[**TOP**]:= **ITEM**. [Insert ITEM in new TOP position]

Step 4. Return



1. Since $TOP=3$, go to Step 2
2. $TOP=3+1=4$
3. $STACK [TOP]=STACK[4]= WWW$
4. Return

STACKS: Popping Algorithm



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POP (STACK, TOP, ITEM)

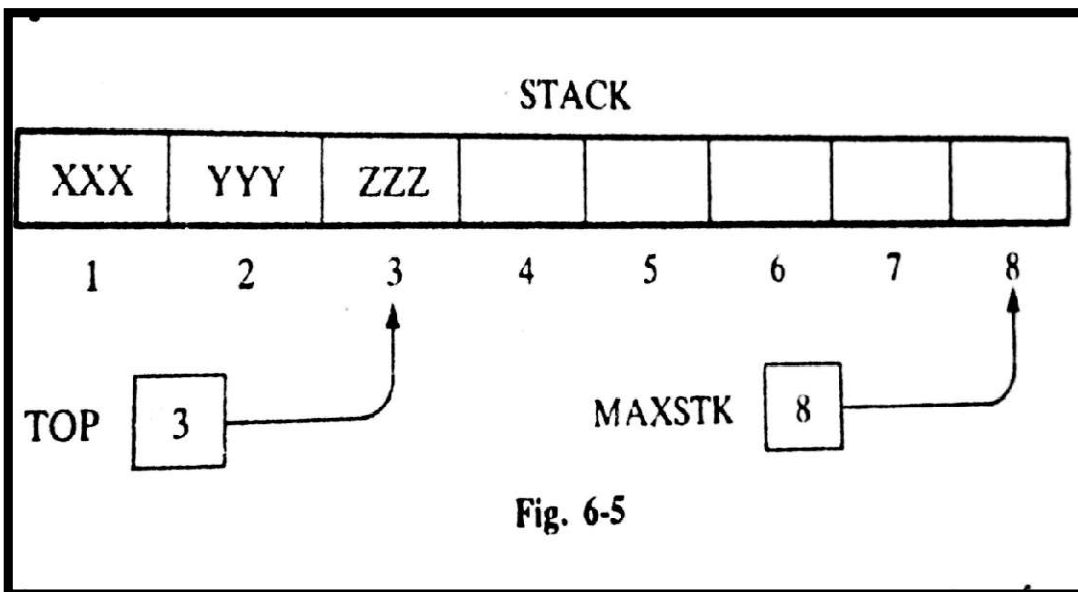
Step 1. [Stack has an item to be removed?]

If **TOP=0**, then print: **Underflow**, and Return.

Step 2. Set **ITEM=STACK[TOP]** [Assign TOP element to ITEM.]

Step 3. Set **TOP:=TOP-1**. [Decreases TOP by 1]

Step 4. Return



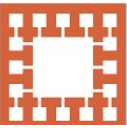
1. Since TOP=3, go to Step 2

2. ITEM=ZZZ.

3. TOP=3-1=2.

4. Return.

STACKS:



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PUSH

1. Since $TOP=3$, go to Step 2
2. $TOP=3+1=4$
3. $STACK[TOP]=STACK[4]= WWW$
4. Return

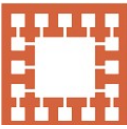
POP

1. Since $TOP=3$, go to Step 2
2. $ITEM=ZZZ.$
3. $TOP=3-1=2.$
4. Return.

✓ **TOP** is changed before the insertion in PUSH, whereas

✓ **TOP** is changed after the deletion in POP

STACKS: Arithmetic Expressions;



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✓ Let Q be an arithmetic expression involving constants and operation.

$Q: 2 \uparrow 3 + 5 * 2 \uparrow 2 - 12 / 6$, we have to find the value of Q

✓ Q may have different levels of precedence in its binary operations.

✓ We assume the following three levels of precedence for the usual **FIVE** binary operations.

➤ Highest: Exponential (\uparrow)

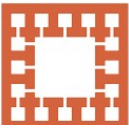
➤ Next Highest: Multiplication ($*$) and division ($/$)

➤ Lowest: Addition ($+$) and Subtraction ($-$)

✓ Thus, we obtain after exponentiations $(8 + 5 * 4 - 12 / 6)$

✓ After, multiplication and division $(8 + 20 - 2)$

✓ After. Addition and subtraction , 26



STACKS: Polish Notation

For most common arithmetic operations:

$A+B$, $C-D$, $E * F$, G/H (*infix notation*)

Polish notation refers to the notation in which the operator symbol is placed before its two operands.

For example: $+AB$, $-CD$, $*EF$, $/GH$

Instant Exercise: (*infix expression* to *polish notation*)

$$(A+B) * C = ?$$

$$= *[A+B]C=?$$

$$= *+ABC$$

$$A+(B * C) = ?$$

$$= A+[*BC]=?$$

$$= +A*BC$$

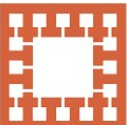
$$(A+B)/(C-D) = ?$$

$$= [+AB]/[-CD] = ?$$

$$= /+AB-CD$$

- ✓ The fundamental property of polish notation is that the order in which the operations are to be performed is completely determined by
 - ✓ The positions of the operators, and
 - ✓ Operands in the expression.
- ✓ **One never needs parentheses when writing expressions in polish notation.**

STACKS: Reverse Polish Notation



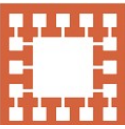
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Reverse Polish Notation refers to the analogous notation in which the operator symbol is placed after its two operands:

AB+, CD-, EF*, GH/

- ✓ This notation is frequently called *postfix notation*
- ✓ One never needs parentheses to determine the order of the operations in any arithmetic expression in *postfix notation*

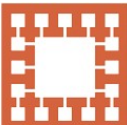
STACKS: Evaluation of Postfix Expression



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P: 5, 6, 2, +, *, 12, 4, /, -

Symbol Scanned	STACK
(1) 5	5
(2) 6	5, 6
(3) 2	5, 6, 2
(4) +	5, 8
(5) *	40
(6) 12	40, 12
(7) 4	40, 12, 4
(8) /	40, 3
(9) -	37
(10))	



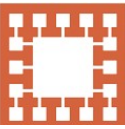
STACKS:

Algorithm to find the VALUE of an arithmetic expression P written in postfix notation.

- 1. Add a right parenthesis “)” at the end of P.**
 - 2. Scan P from left to right and repeat Steps 3 and 4 for each of P until the “)” is encountered.**
 - 3. If an operand is encountered, put it on STACK**
 - 4. If an operator \otimes is encountered, then:**
 - a) Remove the two top elements of the STACK, where A is the top element and B is the next-to-top element**
 - b) Evaluate $B \otimes A$**
 - c) Place the result of (b) back on STACK.**

[End of If structure]
[End of Step 2 loop]
 - 5. Set VALUE equal to the TOP element on STACK.**
- Exit.**

STACKS: Transforming Q into P expression

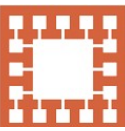


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Q: $A + (B * C - (D / E \uparrow F) * G) * H = (?)$ P expression

Symbol Scanned	Stack	Expression P
(1) A	(A
(2) +	(+	A
(3) ((+ (A
(4) B	(+ (A B
(5) *	(+ (*	A B
(6) C	(+ (*	A B C
(7) -	(+ (-	A B C *
(8) ((+ (- (A B C *
(9) D	(+ (- (A B C * D
(10) /	(+ (- (/	A B C * D
(11) E	(+ (- (/	A B C * D E
(12) ↑	(+ (- (/ ↑	A B C * D E
(13) F	(+ (- (/ ↑	A B C * D E F
(14))	(+ (-	A B C * D E F ↑ /
(15) *	(+ (- *	A B C * D E F ↑ /

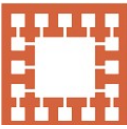
STACKS: Transforming Q into P expression



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Q: $A + (B * C - (D / E \uparrow F) * G) * H = (?)$ P expression

Symbol Scanned	Stack	Expression P
(7) -	(+ (-	A B C *
(8) ((+ (- (A B C *
(9) D	(+ (- (A B C * D
(10) /	(+ (- (/	A B C * D
(11) E	(+ (- (/	A B C * D E
(12) ↑	(+ (- (/ ↑	A B C * D E
(13) F	(+ (- (/ ↑	A B C * D E F
(14))	(+ (-	A B C * D E F ↑ /
(15) *	(+ (- *	A B C * D E F ↑ / ¹⁸
(16) G	(+ (- *	A B C * D E F ↑ / G
(17))	(+	A B C * D E F ↑ / G * -
(18) *	(+ *	A B C * D E F ↑ / G * -
(19) H	(+ *	A B C * D E F ↑ / G * - H
(20))		A B C * D E F ↑ / G * - H * +



Instant TEST

12 , 7 , 3 , - , / , 2 , 1 , 5 , + , * , +

Which Expression?

Postfix or Prefix

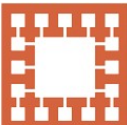
✓ **Postfix**

Equivalent infix expression ?

$$\begin{aligned} P &= 12, [7-3], /, 2, 1, 5, +, *, + \\ &= [12/(7-3)], 2, 1, 5, +, *, + \\ &= [12/(7-3)], 2, [1+5], *, + \\ &= [12/(7-3)], [2 * (1+5)], + \\ &= 12/(7-3) + 2 * (1+5) \end{aligned}$$

Result: 15

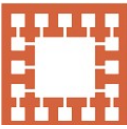
Instant TEST



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Infix	Prefix	Postfix
$A+B * C+D$	$++A*BCD$	$ABC*+D+$
$(A+B) * (C+D)$	$*+AB+CD$	$AB+CD+*$
$A*B + C*D$	$+*AB*CD$	$AB*CD*+$
$A+B+C+D$	$++AB+CD$	$AB+CD++$

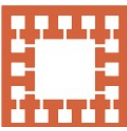
Instant TEST



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Infix Expression	Prefix Expression	Postfix Expression
$1-4^3+7*(9^1/5)-2$	$-1+^43-*7/^9152$	$143^--791^5/*+2-$
$A+B-(C*D^E)*X+Y$	$+A-B+**C^DEXY$	$AB+CDE^*X*-Y+$

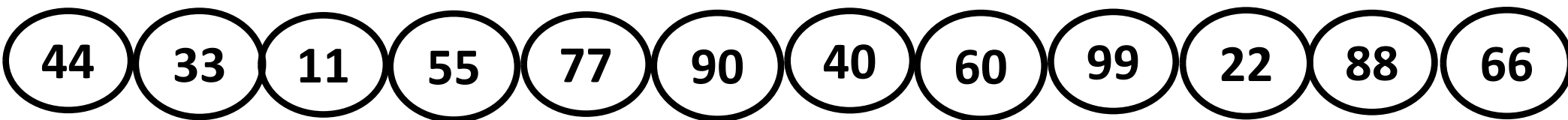
QUICKSORT: An Application of STACKS



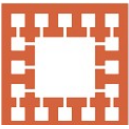
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- What is Sorting?
 - ✓ The operation of rearranging the elements of a list so that they are in some logical order.
 - Numerically ordered (When list contain numerical data)
 - Alphabetically ordered (When list contains character data)
- Quicksort is an algorithm of the **DIVIDE-AND-CONQUER** type.
 - The problem of sorting a set is reduced to the problem of sorting two smaller sets.
- We illustrate this “Reduced Step” by means of a specific example
- Suppose A is the following list of 12 numbers.

A

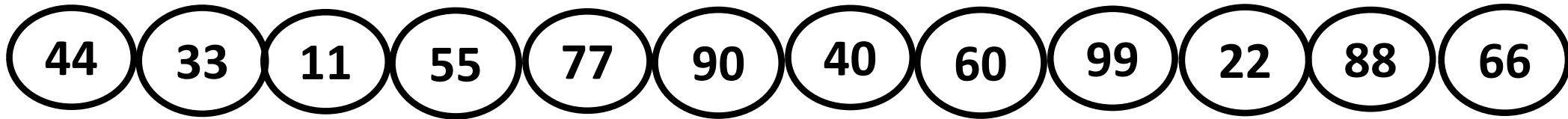


Reduction Step of the Quicksort Algorithm



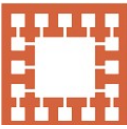
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A

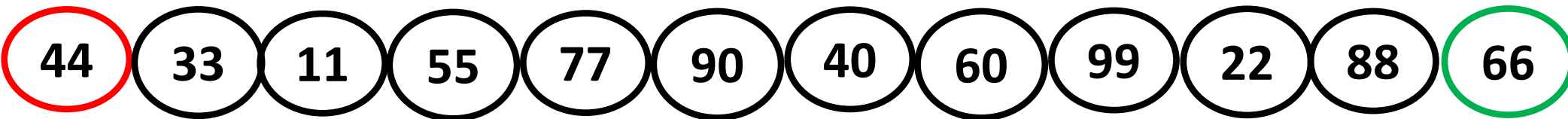
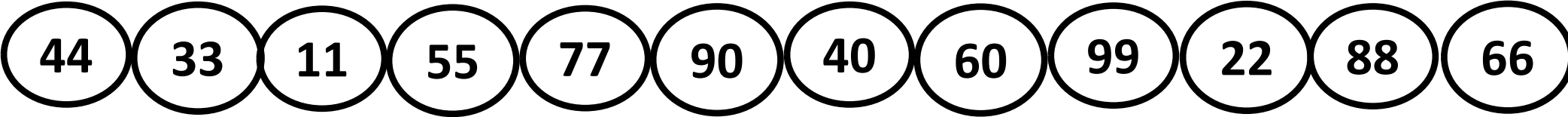


- The reduction step of the quicksort algorithm finds the final position of one of the numbers.
- In this illustration, we use the first number 44.

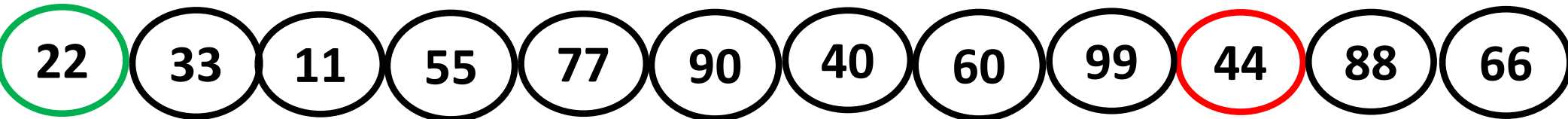
STACKS: Quicksort/ Reduction Step



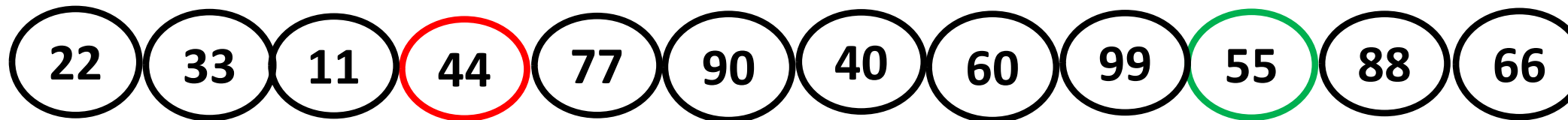
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- ✓ Beginning with the last number, **66**, scan the list from right to left till less than **44**.
- ✓ Interchange **44** and **22**

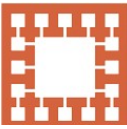


- ✓ Beginning with, **22**, scan left to right till greater than **44**
- ✓ Interchange **44** and **55**

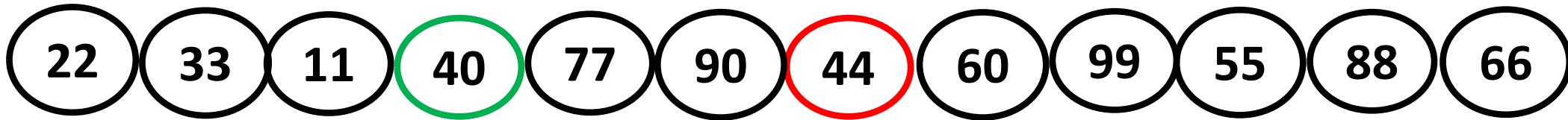


- ✓ Beginning with, **55**, scan the list from right to left till less than **44**.
- ✓ Interchange **44** and **40**

STACKS: Quicksort/ Reduction Step

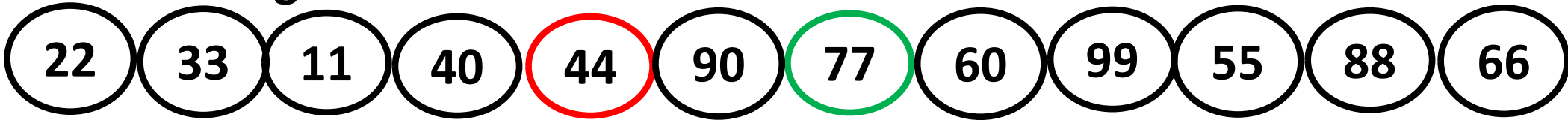


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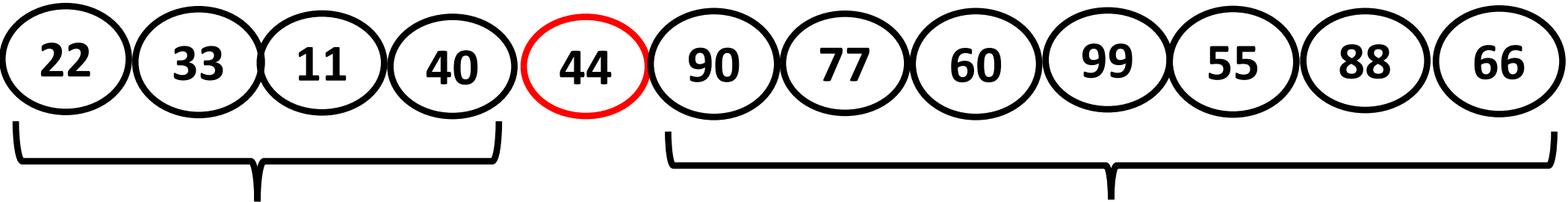
✓ Beginning with, 40, scan left to right till greater than 44

✓ Interchange 44 and 77



✓ Beginning with, 77, scan the list from right to left till less than 44.

✓ Do not meet such a number before meeting 44.



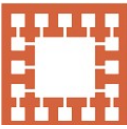
First sublist

Second sublist

✓ Thus, 44 is correctly placed in its final position.

✓ The task of sorting the original list A has now been reduced to the task of sorting each of the above sublists.

STACKS: Quicksort/ Reduction Step



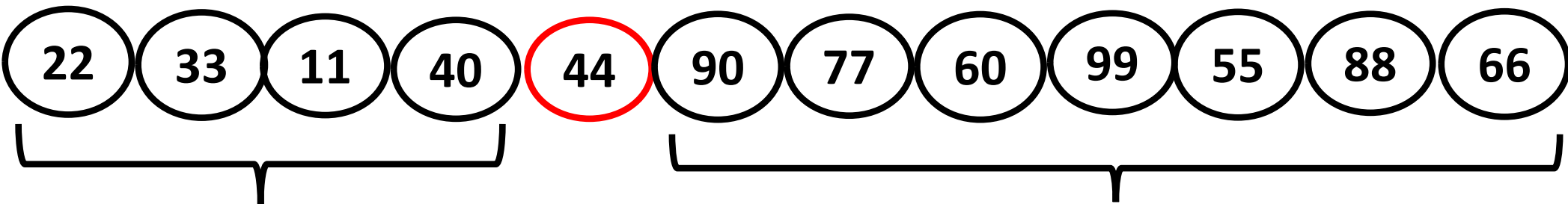
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- The reduction step is repeated with each sublist containing 2 or more elements.
- Since we can process one sublist at a time, we must be able to keep track of some sublist for future processing.
- **This is accomplished by using two STACKS**

LOWER and UPPER

to temporarily hold such sub-lists.

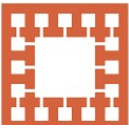
- The addresses of the first and last elements of each sublist, called its “**BOUNDARY VALUES**”, are pushed onto the STACKs **LOWER** and **UPPER**, respectively.
- The reduction steps is applied to a sublist only after its **BOUNDARY VALUES** are removed from the STACKs.



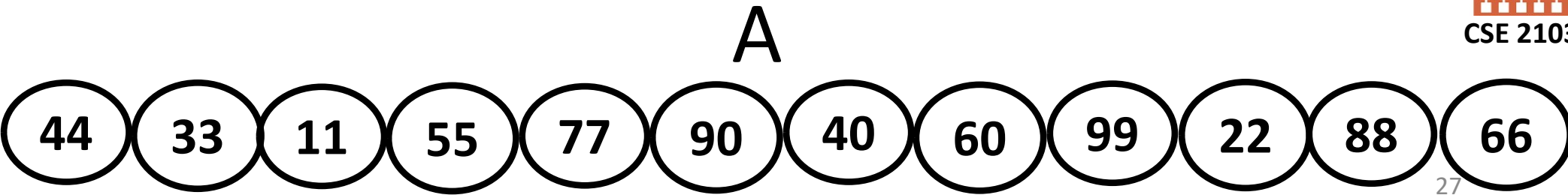
First sublist

Second sublist

STACKS: Illustration of the way the STACKS LOWER and UPPER are used



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N=12 elements,

Thus,

Boundary values are ()?

1 and 12.

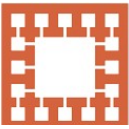
Now,

1 and 12 should be Stacked

LOWER:1 and **UPPER:12**

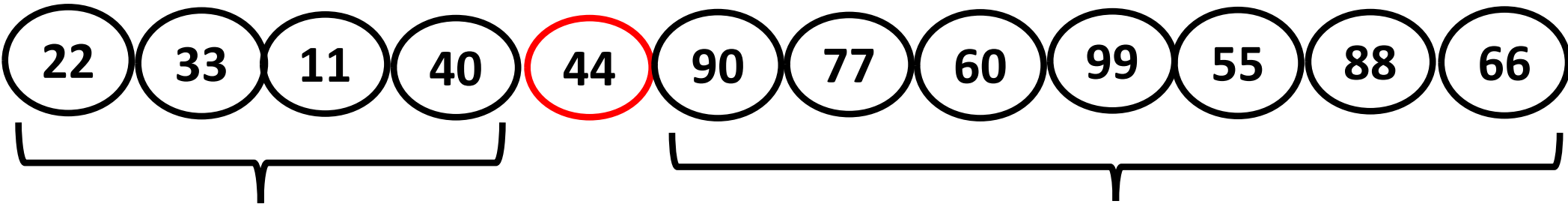
- In order to apply the **REDUCTION STEP**, the algorithm first removes the top values 1 and 12.
- After removing the top values 1 and 12 from the STACK, leaving **LOWER:(Empty)** and **UPPER: (Empty)**
- Then Applies the **REDUCTION STEP** to the corresponding list A[1], A[2],...,A[12].

STACKS: Illustration of the way the STACKS LOWER and UPPER are used



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- After executing **REDUCTION STEP** to the list A[1] to A[12]
 - Finally places the first element 44, in A[5].



First sublist

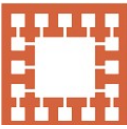
Second sublist

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- Accordingly, the algorithm pushes the boundary values
 - 1 and 4 of the first sublist, and
 - 6 and 12 of the second sublist on to the STACK to yield
LOWER= **1, 6** and UPPER= **4, 12**
- In order to apply the **REDUCTION STEP** again, the algorithm removes the TOP values 6 and 12 from the STACKs, leaving
LOWER= **1**, and UPPER= **4**

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STACKS: Illustration of the way the STACKS LOWER and UPPER are used



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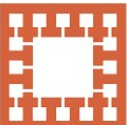
A[6]	A[7]	A[8]	A[9]	A[10]	A[11]	A[12]
90	77	60	99	55	88	66
66	77	60	99	55	88	90
66	77	60	90	55	88	99
66	77	60	88	55	90	99

First sublist

Second sublist

- The second sublist has only one element, Accordingly
- The algorithm pushes only the boundary values 6 and 10 of the first sublist on the STACKs to yield
 - LOWER= 1, 6 and UPPER= 4, 10

QUICKSORT



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- The quick sort is regarded as the best sorting algorithm.
- This is because of its significant advantage in terms of efficiency because it is able to deal well with a huge list of items.
- Because it sorts in place, no additional storage is required as well.
- The slight disadvantage of quick sort is that its worst-case performance is similar to average performances of the bubble, insertion or selections sorts.
- In general, the quick sort produces the most effective and widely used method of sorting a list of any item size.