

# **Department Of Computer Science and Engineering**

**Course Title:** Operating System Lab

**Course Code:** CSE 406

**Title:** CPU Scheduling Priority Scheduling Algorithm

Submitted To Submitted By

Atia Rahman Orthi Md. Atikul Islam Atik
Lecturer Reg No:21201063
Department Of Computer Science & Sec:B1

Engineering

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# **Problem:**

Process	Priority	Arrive Time	Burst Time	
1	10 0		5	
2	20	1	4	
3	30	2	2	
4	40	4	1	

# **Output:**

Process	Priority	Arrival	Burst	СТ	TAT	WT	RT
1	10	0	5	12	12	7	0
2	20	1	4	8	7	3	0
3	30	2	2	4	2	0	0
4	40	4	1	5	1	0	0

### **Source Code:**

```
• • •
            time = 0
gantt_chart
             completed = {}
backup = {}
             backup =
             first_response = {}
             \begin{tabular}{ll} \# \ Sort \ processes \ by \ arrival \ time \ initially \\ process\_list.sort(key=lambda \ x: \ x[3]) \end{tabular} 
             for process in process_list:
   pid = process[1]
   backup[pid] = {"arrival_time": process[3], "burst_time": process[2], "priority": process[0])
   first response[oid] = -1
             remaining_processes = process_list[:]
            while remaining_processes:
                    available_processes
                    for p in remaining_processes:
    if p[3] <= time:
        available_processes.append(p)</pre>
                   if not available_processes:
    gantt_chart.append("Idle")
    time += 1
    continue
                   # Select the highest-priority process (higher number = higher priority) process = \max(available\_processes, key=lambda x: x[0])
                    # Record first response time if it's the first execution
if first_response[pid] == -1:
    first_response[pid] = time - process[3]
                   execution_time = min(time_quanta, process[2])
                   gantt_chart.append(pid)
                   time += execution_time
process[2] -= execution_time
                   # If process is completed
if process[2] == 0:
                          process[2] == 0:
completion_time = time
arrival_time = backup[pid]["arrival_time"]
burst_time = backup[pid]["burst_time"]
turnaround_time = completion_time - arrival_time
waiting_time = turnaround_time - burst_time
response_time = first_response[pid]
                          {\tt remaining\_processes.remove(process)}
             return {"gantt_chart": gantt_chart, "completed": completed}
     def print_table(completed):
    print("\nProcess\tArrival Time\tBurst Time\tPriority\tCompletion Time\tTurnaround Time\tWaiting Time\tResponse Time")
             for key, values in completed.items():
    print(f"{key}\t{values[0]}\t\t{values[1]}\t\t{values[2]}\t\t{values[3]}\t\t{values[4]}\t\t{values[5]}\t\t{values[6]}")
     def collect_input():
    process_list = []
    n = int(input("Enter the number of processes: "))
    time_quanta = int(input("Enter the time quanta: "))  # Collecting time quanta
                   l in range(n):
priority = int(input(f'Enter the priority for process {i+1} (Higher value -> Higher priority): '))
pid = f"P{i+1}"
arrival_time = int(input(f'Enter the arrival time for process {i+1}: '))
burst_time = int(input(f'Enter the burst time for process {i+1}: '))
process_list.append([priority, pid, burst_time, arrival_time])
             return process_list, time_quanta
            __name__ == "__main__":
inputs = collect_input()
result = priority_preemptive(inputs[0], inputs[1])
            print("\nGantt Chart:", result["gantt_chart"])
print_table(result["completed"])
```

# **Output:**

```
■os/lab4 □ main !? □ v3.13.2 □ 09:54 PM
 python -u "/home/atik/Codes/python/os/lab4/priority_premptive.py"
Enter the number of processes: 4
Enter the time quanta: 1
Enter the priority for process 1 (Higher value -> Higher priority): 10
Enter the arrival time for process 1: 0
Enter the burst time for process 1: 5
Enter the priority for process 2 (Higher value -> Higher priority): 20
Enter the arrival time for process 2: 1
Enter the burst time for process 2: 4
Enter the priority for process 3 (Higher value -> Higher priority): 30
Enter the arrival time for process 3: 2
Enter the burst time for process 3: 2
Enter the priority for process 4 (Higher value -> Higher priority): 40
Enter the arrival time for process 4: 4
Enter the burst time for process 4: 1
Gantt Chart: ['P1', 'P2', 'P3', 'P3', 'P4', 'P2', 'P2', 'P2', 'P1', 'P1', 'P1', 'P1']
Process Arrival Time Burst Time
                                                     Completion Time Turnaround Time Waiting Time Response Time
                                      Priority
                                      40
                                      20
                                      10
  atik ■os/lab4 □ main !? □ v3.13.2 □ 09:55 PM
```

# Algorithm:

**Input Process Details**: Each process should have an ID, priority, arrival time, and burst time.

#### **Initialize Variables:**

- Set currentTime = 0.
- Set remainingTime = burstTime for all processes.

#### Find the Next Process to Execute:

- Identify the highest-priority process that has already arrived and still has remaining execution time.
- If multiple processes have the same priority, use arrival time as a tiebreaker (earlier arrival gets preference).

#### **Execute the Process:**

- Run the selected process for 1 time unit.
- Decrease its remaining Time by 1.

## **Check for Process Completion:**

- If a process finishes execution (remainingTime == 0), record:
  - Completion Time (CT) = currentTime + 1.
  - o **Turnaround Time (TAT)** = CT Arrival Time.
  - o **Waiting Time (WT)** = TAT Burst Time.

Repeat Steps 3-5 until all processes are completed.

### **Display Results:**

- Print **CT, TAT, WT** for each process.
- Optionally, calculate and display average TAT and WT for overall evaluation.

## **Conclusion:**

- Priority scheduling is one of the most common scheduling algorithms used by the operating system to schedule processes based on their priority. Each process is assigned a priority.
  - **Waiting Time (WT):** The time a process spends waiting in the queue before it starts executing.
  - **Response Time (RT):** The time from process arrival to its first execution.

The output clearly shows the execution order in the Gantt chart and provides detailed scheduling results for each process. The results demonstrate that the Round Robin algorithm ensures fair CPU allocation by giving each process a fixed time slice, reducing starvation. However, selecting an appropriate time quantum is crucial to balancing responsiveness and efficiency.