

# QUEEN MARY, UNIVERSITY OF LONDON

MTH6150

## Numerical Computing in C and C++

### Exercise Sheet 6

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1. Write a function `cube` that returns the cube of a `double` variable (with return type `double`). Write another function of the same name that returns the cube of an `int` variable (with return type `int`).

Use both of them in the same program, first with the functions declared and defined before `main`, and then with the function bodies (the definition) after `main`.

Modify the `int` version to the following:

```
double cube(const int x){
```

What happens now if you try to run a program with this version and the following version at the same time:

```
double cube(const double x){ ?
```

2. Type in the function for the factorial on slide 4 of lecture 6, and run it to find out if it is correct. Add a command within the function to output `n` to the screen and run this code to verify what sequence of numbers the function is called with.
3. Run the example pieces of code on slides 6 to 10 to understand which part of the code variables exist and why the code won't compile in some cases.
4. We have met the function `pow` in `<cmath>`, with  $z = x^y$  given by

```
double z = pow(x, y);
```

When `y` is an integer, it may be more efficient to start with 1 and then repeatedly multiply by `x`, doing this `y` times. Implement a function

```
double pow_int(const double x, const int k){  
    ...  
}
```

that calculates  $x^k$  using a simple `for` loop. Check the numerical results by comparing them with `pow`, and/or with examples where you know the answer already.

5. In question 4, the number of multiplications can be reduced by noting that

$$x^{2m} = (x^m)^2$$

$$x^{2m+1} = x (x^m)^2$$

Implement a variant of `pow_int` that takes advantage of this.

One way would be using recursive function calls:

```
double pow_int_rec(const double x, const int k){
    if(k==1){
        ...
    }
    else if(k>=2 && k%2==0){
        const double w = pow_int_rec(x, k/2);
        return w*w;
    }
    else{
        ...
    }
}
```

As before, you can check the numerical results by comparing with `pow`.

Can you see how the computing time to calculate  $x^k$  increases as  $k$  increases for each of the two methods (in theory, not by trying to time the code)?

6. The probability density function (pdf) of the normal distribution with mean  $\mu$  and standard deviation  $\sigma$  is given by

$$\varphi(x; \mu, \sigma) = \frac{1}{\sigma\sqrt{2\pi}} e^{-\frac{(x-\mu)^2}{2\sigma^2}}$$

Write a function `norm_pdf` with three arguments  $(x, \mu, \sigma)$  that returns this pdf.

```
double norm_pdf(const double x, const double mu, const double sigma){
    ...
}
```

As a check, two example values are  $\varphi(1; 0, 1) \approx 0.242$ ,  $\varphi(2; 1, 2) \approx 0.176$ .

7. The standard normal cumulative distribution function (i.e. for  $\mu = 0$ ,  $\sigma = 1$ )

$$\Phi(x) = \int_{-\infty}^x \varphi(t; 0, 1) dt = \int_{-\infty}^x \frac{1}{\sqrt{2\pi}} e^{-t^2/2} dt$$

does not have an explicit closed-form expression. It can be expressed in terms of an infinite series. As an intermediate step

$$\Phi(x) = \frac{1}{2} [1 + g(x / \sqrt{2})]$$

where  $g(w)$  is known as the error function. We also have

$$g(w) = \frac{2}{\sqrt{\pi}} \sum_{k=0}^{\infty} \frac{(-1)^k w^{2k+1}}{k!(2k+1)}$$

Write a function that calculates an approximation to the error function  $g(w)$  using the first  $n$  terms from this infinite series:

```
double erff(const double w, const int n){
    ...
}
```

Then write a function

```
double norm_cdf(const double x, const int n){
    ...
}
```

that calculates  $\Phi(x)$  using the formula for  $\Phi$  in terms of  $g$ , i.e. by using the function `erff` inside the function `norm_cdf`.

To check the results, some standard values that are used for calculation of 90% and 95% confidence intervals are  $\Phi(1.64) \approx 0.95$  and  $\Phi(1.96) \approx 0.975$ . Also,  $\Phi(x) = 1 - \Phi(-x)$  for all real  $x$ .

Note that the series is very slow to converge for values of  $x$  outside the range  $-2.5$  to  $2.5$ , so larger values of  $n$  are needed.  $n = 10$  or  $15$  is reasonable inside this range.

Also, see if you can make the function `erff` more efficient by avoiding calculating  $w^{2k+1}$  and  $k!$  at each loop iteration (by storing certain variables declared before the loop and updating them at each loop iteration).

8. A definite integral can be approximated by the composite trapezoidal rule

$$I = \int_{t_0}^{t_N} f(t)dt \simeq \Delta t \left[ \frac{f(t_0) + f(t_N)}{2} + \frac{f(t_1) + f(t_2)}{2} + \dots + \frac{f(t_{N-1}) + f(t_N)}{2} \right]$$

$$= \Delta t \left[ \frac{f(t_0) + f(t_N)}{2} + \sum_{k=1}^{N-1} f(t_0 + k\Delta t) \right], \quad \Delta t = \frac{t_N - t_0}{N} > 0$$

on a set of equidistant points  $\{t_0, t_1, t_2, \dots, t_N\}$  for some positive integer  $N$ . Use this method to approximate  $\Phi(x)$  from question 7, by taking  $f(x)$  to be  $\varphi(x; 0, 1)$  from question 6. Try this for various values of  $N$ .

Note that  $\Phi(0) = 0.5$ , so for  $x > 0$

$$\Phi(x) = 0.5 + \int_0^x \varphi(t; 0, 1)dt$$

Also,  $\Phi(x) = 1 - \Phi(-x)$  for all real  $x$ .

## Exercise Sheet 6

### Hints for question 7

- Some values for the error function, denoted by  $g(w)$  in the question, are:  
 $g(1) \approx 0.8427008$   
 $g(0.5) \approx 0.5204999$   
Also,  $g(-w) = -g(w)$  for all real  $w$ .
- Using a loop to calculate a sum: there are simpler examples on exercise sheet 4, question 1b or question 3 (code as a solution to question 3 is on QMplus).
- There is code for the factorial  $n!$  on slide 4 of lecture 6.
- In the calculations, possible sources of errors are:  
integer division, so convert `int` to `double`;  
and use round brackets to ensure that the correct terms are included in the denominator of a fraction.