

HASHING WITH CHAINING

Dictionary: Abstract Data Type (ADT)
maintain set of items, each
with a key

$O(n)$ {
via
AVL
↓
 $O(1)$

- insert(item) (overwrite any existing key)
- delete(item)
- Search(key): return item with
given key
or report doesn't
exist

Python: dict

$D[key]$ ~ search
 $D[key] = val$ ~ insert
del $D[key]$ ~ delete
 $\rightarrow item = (key, value)$

Motivation:

- document
- database
- compilers & interpreters
- network router
server
- substring search

- string commonalities
- file synchronization
- cryptography

Simple approach:

Direct-access table

1	/
2	item
3	/
	..
	..
	..
	item
	/

Badness:

- 1) keys may not be nonneg. integers
- 2) gigantic memory hop

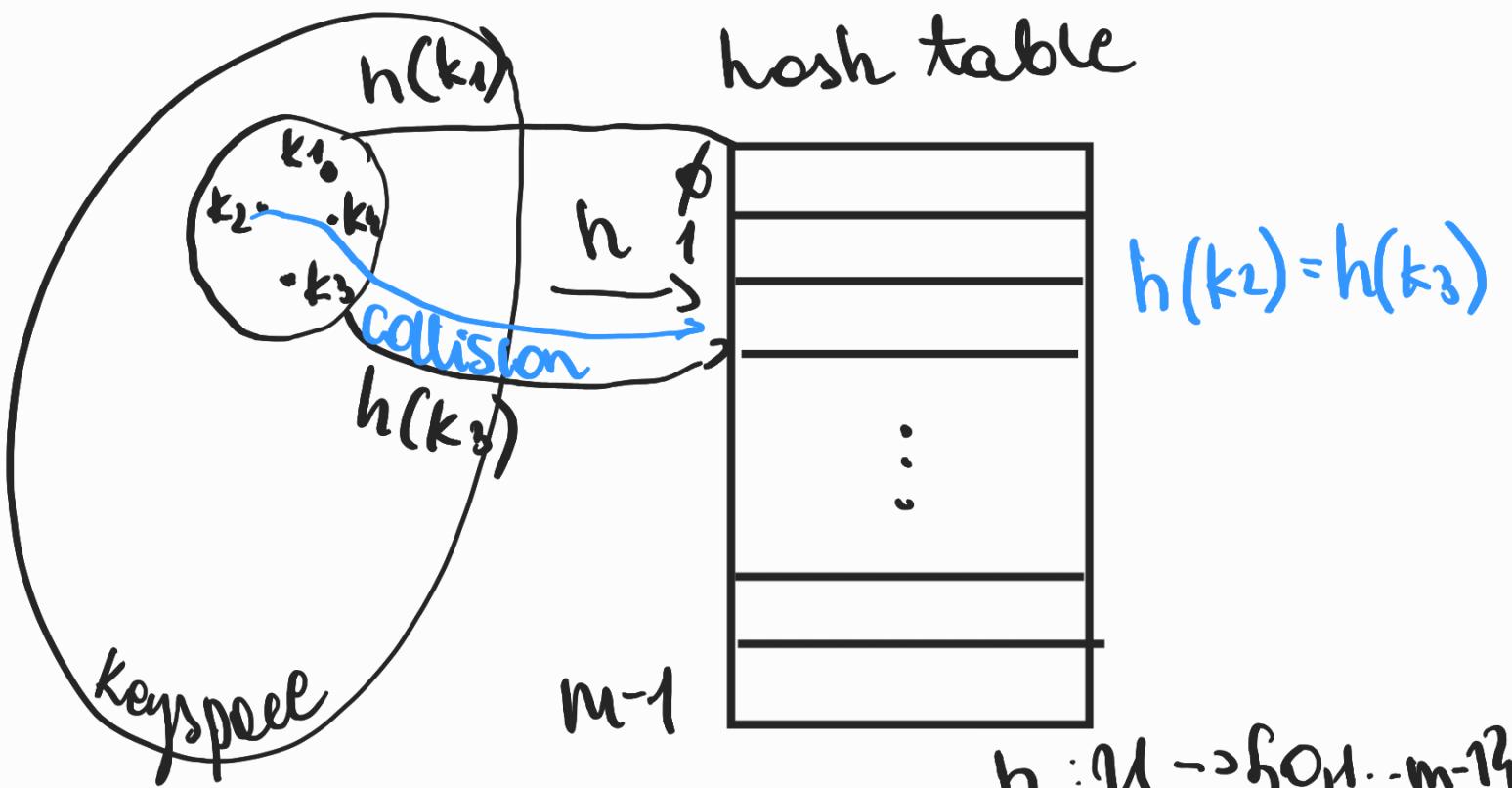
Solution to 1: prehash

- maps keys to nonneg. integers
- in theory keys are finite & discrete (string of bits)

- in Python `hash(x)` is the prehash of x
- $\text{hash('1\phiB')} = \text{hash('1\phiC')} = 64$
ideally: $\text{hash}(x) = \text{hash}(y) \Leftrightarrow x = y$

Solution to 2: hashing

- reduce universe \mathcal{U} of all keys (integers) down to reasonable size m for table



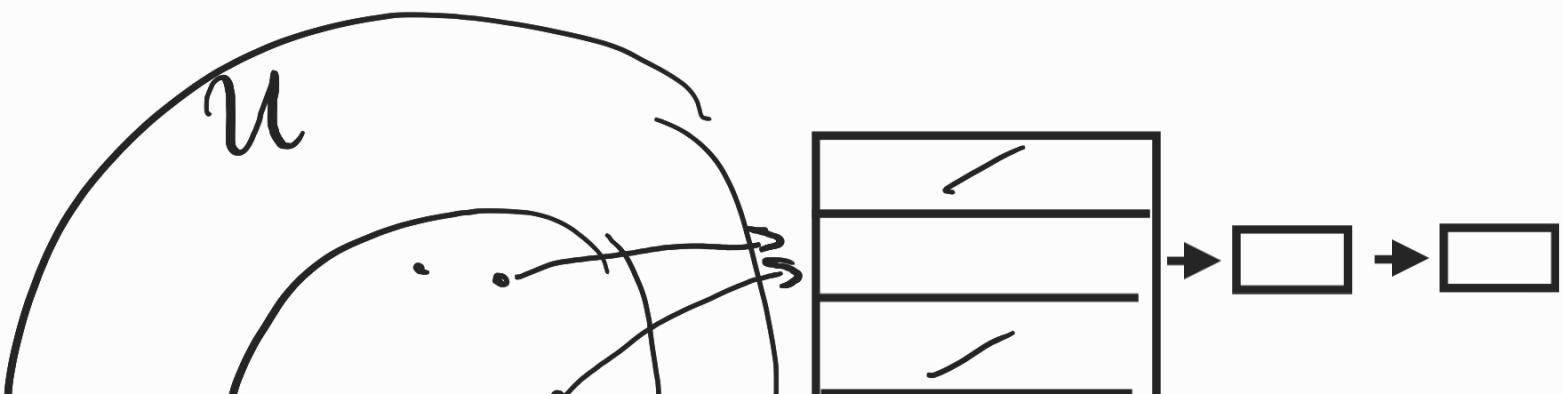
- idée : $m = \Theta(n)$

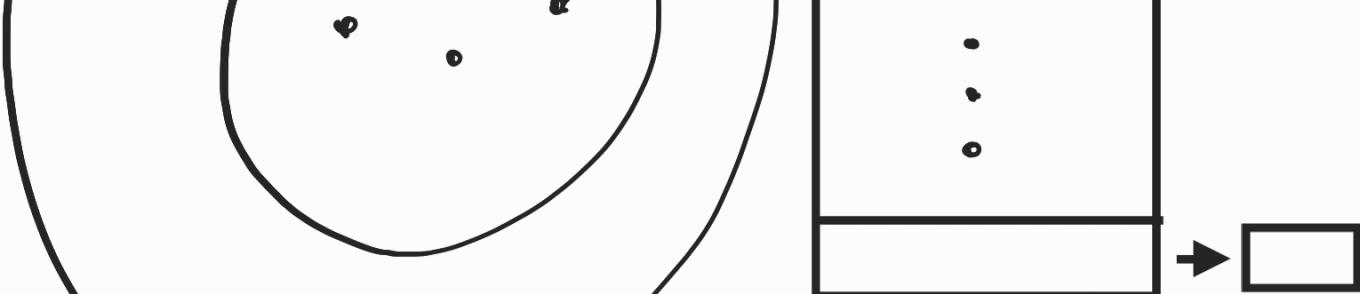
↓
keys in dictionary

DEALING WITH COLLISIONS

CHAINING:

linked list of colliding elements in each slot of hash table





Worst case: $\Theta(n)$ ← if we are unlucky with hash function h

Simple uniform hashing:

each key is equally likely to be hashed to any slot of the table, dependent of where other keys hashing.

Analysis:

- expected length of chain for n keys and m slots
 $= \underbrace{\frac{1}{m} + \frac{1}{m} + \dots + \frac{1}{m}}_n = \frac{n}{m}$

$$\frac{n}{m} = d = \frac{\text{load}}{\text{factor}}$$

$$= \Theta(1) \text{ if } m = \Theta(n)$$

$$\Rightarrow \text{running time} = O(1 + (\text{chain})) \\ = O(1 + \lambda)$$

Hash functions:

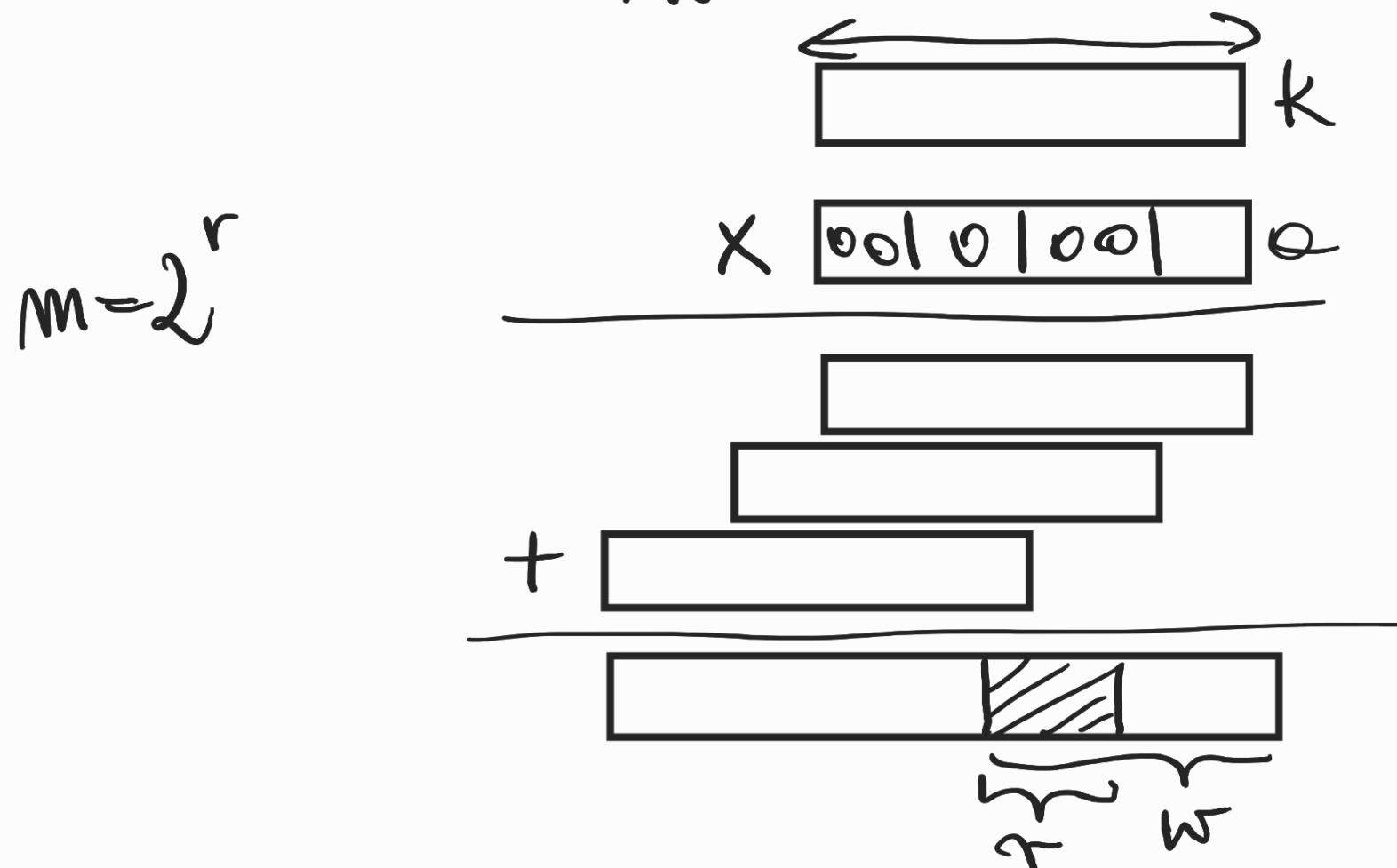
① division method:

$$h(k) = k \bmod m$$

② multiplication method:

$$h(k) = [(a \cdot k) \bmod 2^w] \gg (w-r)$$

$\hookrightarrow w \text{ bits}$



③ Universal hashing:

$$h(k) = [(ck + b) \bmod p] \bmod m$$

↙ random
 $\in \{0, \dots, p-1\}$ ↙ prime
 $> |U|$

for worst-case keys $k_1 \neq k_2$:

$$\Pr[h(k_1) = h(k_2)] = \frac{1}{m}$$

