

ALGORITHMS THINKING, PEAK FINDING

PEAK FINDING

One-dimensional

a	b	c	d	e	f	g	h	i
1	2	3	4	5	6	7	8	9

a - i are numbers

- Position 2 is a peak if and only if $b \geq a$ & $b \geq c$
- Position 9 is a peak if $i \geq h$

PROBLEM: Find a peak if it exists.

Straightforward algorithm

Start from left $\frac{n}{2}$

n

		
--	--	-----	--	--	-----	--

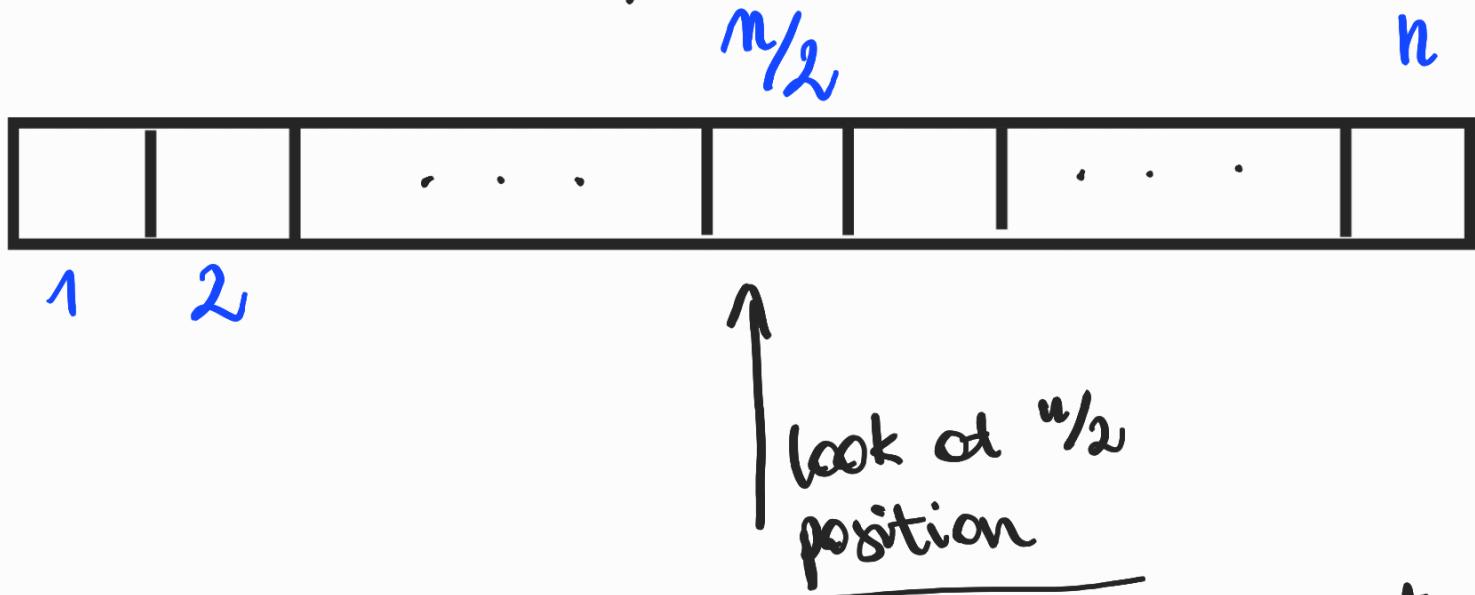
1 2



look at $\frac{n}{2}$ elements.

Worst-case complexity $\Theta(n)$

Divide & conquer



if $a[n/2] < a[n/2-1]$ then only look at left half ($1 \dots n/2-1$) to look a peek.

Else if $a[n/2] < a[n/2+1]$ then ...

$n/2+1 \dots n$ for a peek

Else $n/2$ position is a peek

(Argument that algorithm is correct)

QUESTION?

1 2 3 4 5 6 7 8 9 10 11

$n_1, s_1 \dots n_m, s_m$

$$T(n) = T(n/2) + \Theta(1)$$

↑
work algo
does on
input of
size n

base case: $T(1) = O(1)$

$$T(n) = \underbrace{\Theta(1) + \Theta(1) + \dots + \Theta(1)}_{\log_2 n \text{ times}} = \Theta(\log n)$$

2D VERSION

↑
 n rows
↓

		c	
b	a	d	
	e		

Peak is
when
 $a \geq b, a \geq d,$
 $a \geq c, a \geq e$

← m columns

Greedy Ascent algorithm:

14	15	12	
15	9	11	17
16	17	18	20

$\Theta(mn)$ complexity , when
 $m=n$ then $\Theta(n^2)$
 $j=m/2$

i

Pick middle column $j=\frac{m}{2}$
Find a 1D-peak at (i,j)

Use (i, j) as a start to find a 1D-peak on row i

~~INCORRECT~~

Problem is: 2D peak may not exist on row i

Attempt #2

Pick middle column $j = m/2$
Find global max on column
 j at (i, j) compare $(i, j-1)$,
 (i, j) , $(i, j+1)$.

Pick left columns if
 $(i, j-1) > (i, j)$ (similarly for right).

If (i, j) $\geq (i, j-1), (i, j+1) \Rightarrow$
 $\Rightarrow (i, j)$ is 2-D peak

Solve the new problem with
half the numbers of columns.

When you have a single
column, find the global
max \leftarrow done.

$$T(u, m) = T\left(u, \frac{m}{2}\right) + \Theta(u)$$

$$T(u, 1) = \Theta(u)$$

$$T(u, m) = \underbrace{\Theta(u) + \dots + \Theta(u)}_{\log_2 n \text{ times}}$$

↑
max

$\log_2 n$ times

$$= \Theta(n \log_2 n)$$

