
introduction to AspectJ

CS 119

here viewed as a:

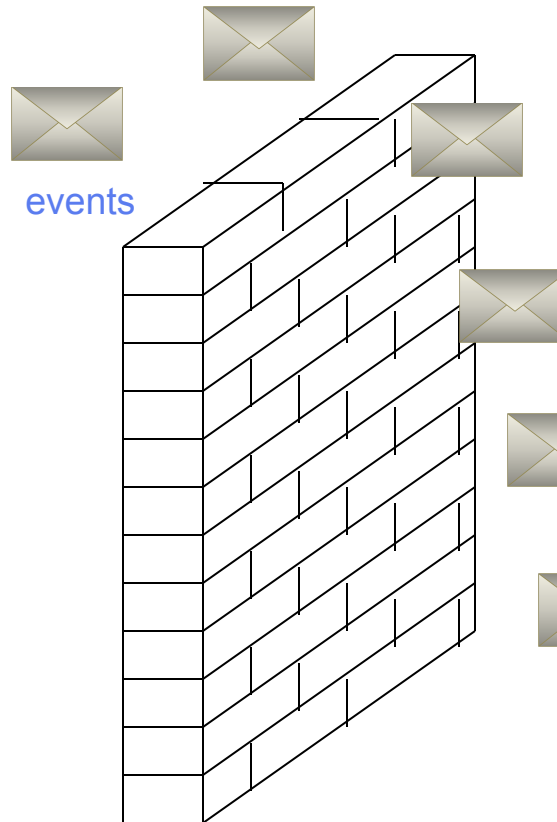
- program instrumentation and
- monitoring framework

monitoring

event generation



instrumentation



event evaluation



specification

why AspectJ?

- so, ... monitoring a program's execution requires these two elements:
 - instrumentation
 - specification
- both elements are provided by **AspectJ**:
 - instrumentation
 - AspectJ's extension to Java
 - specification
 - Java

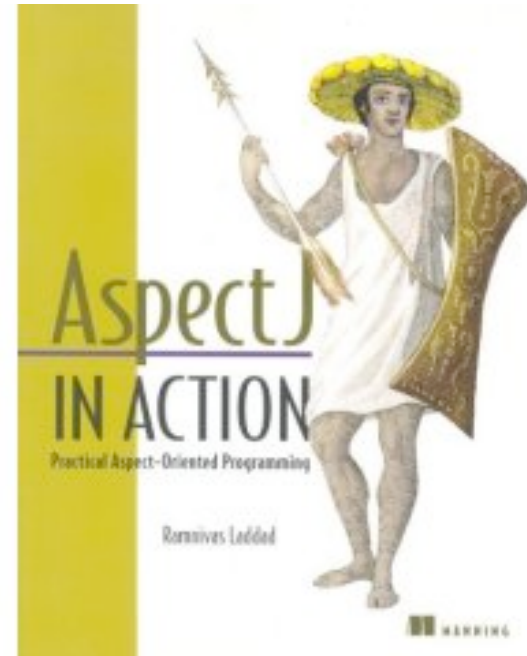
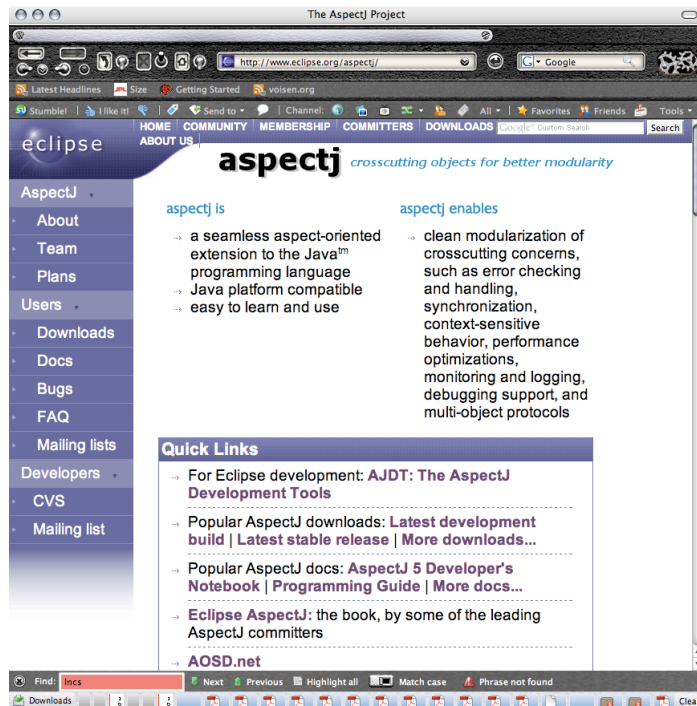
outline

- **this lesson** : introducing the language
- **next lesson** : monitoring with AspectJ

resources

- <http://www.eclipse.org/aspectj>

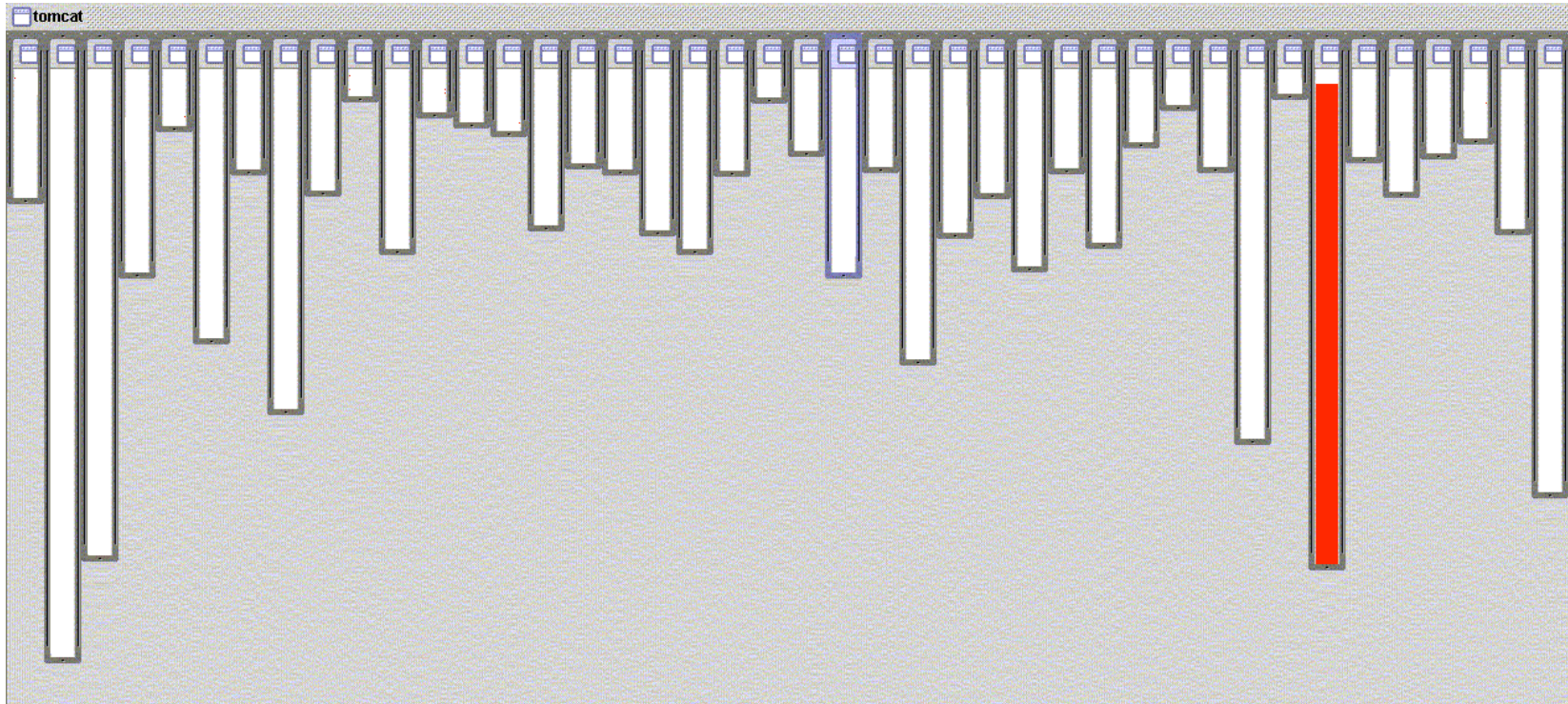
optional reading



AspectJ

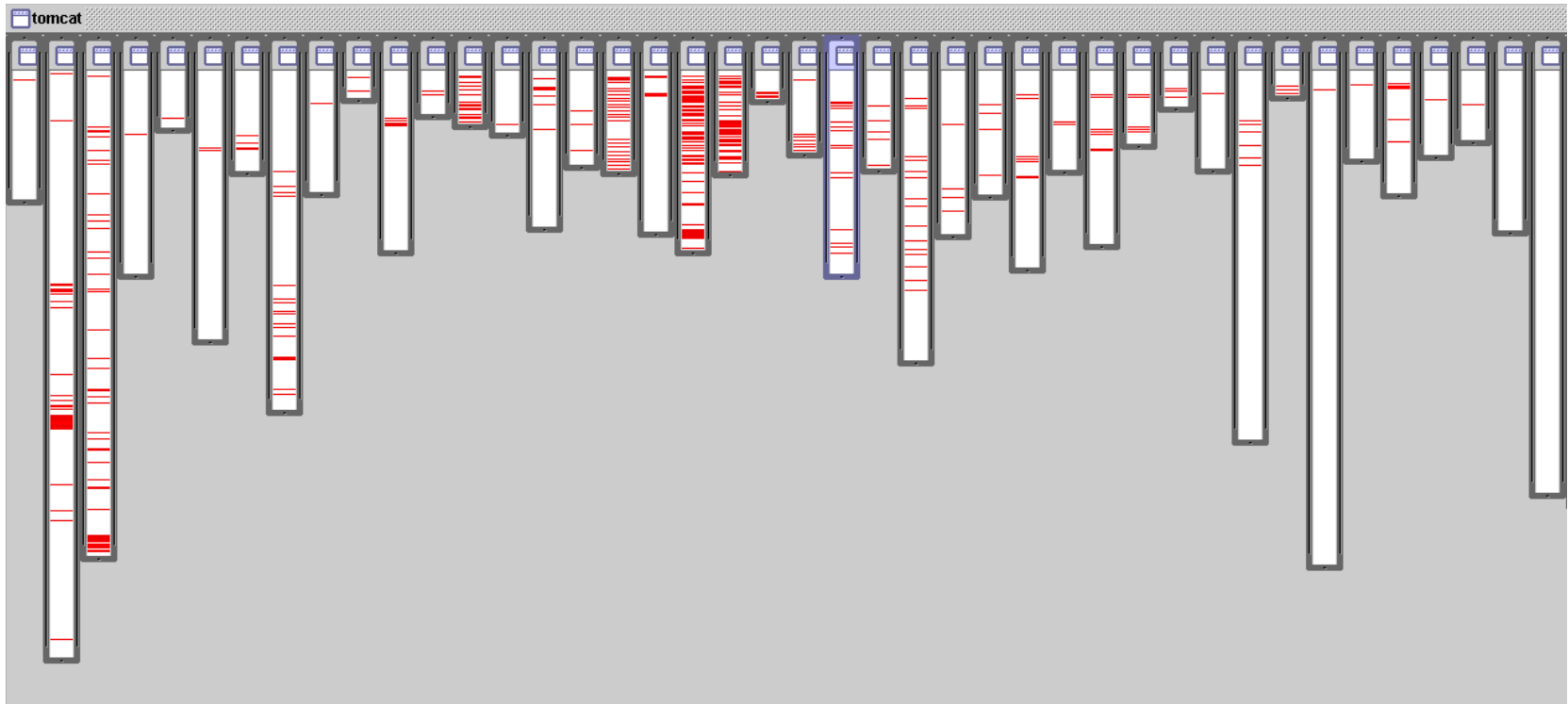
- AspectJ, launched 1998 at Xerox PARC
- an extension of Java
- a new way of modularizing programs compared to object oriented programming
- emphasis on separating out cross-cutting concerns. Logging for example is a concern. That is, code for one *aspect* of the program is collected together in one place
- we shall use it purely for monitoring, and we do not focus on the broader application of AOP as a programming paradigm
- we will, however, briefly explain the more general purpose of AOP
- the AspectJ compiler is free and open source, very mature
- AspectJ works with Eclipse, and other IDEs
- outputs .class files compatible with any JVM

good modularity



- XML parsing in org.apache.tomcat
 - red shows relevant lines of code
 - nicely fits in one box

bad modularity

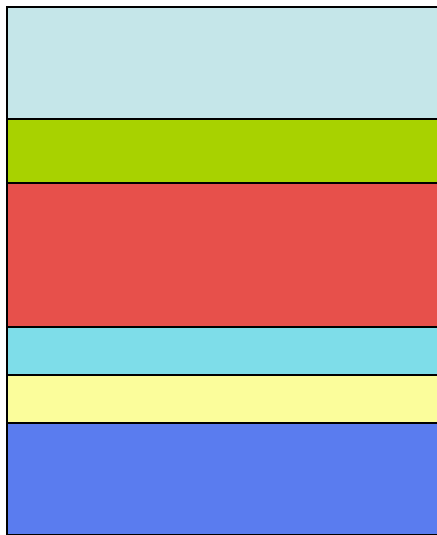


- where is logging in org.apache.tomcat
 - red shows lines of code that handle logging
 - not in just one place
 - not even in a small number of places

two central problems

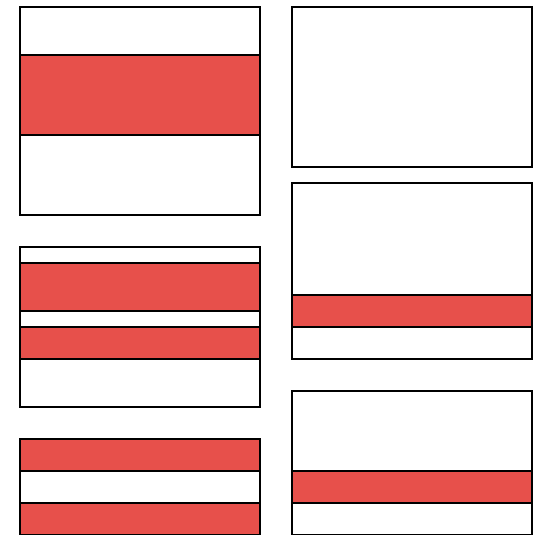
AOP tries to solve

code tangling:
one module
many concerns



example:
logging

code scattering:
one concern
many modules

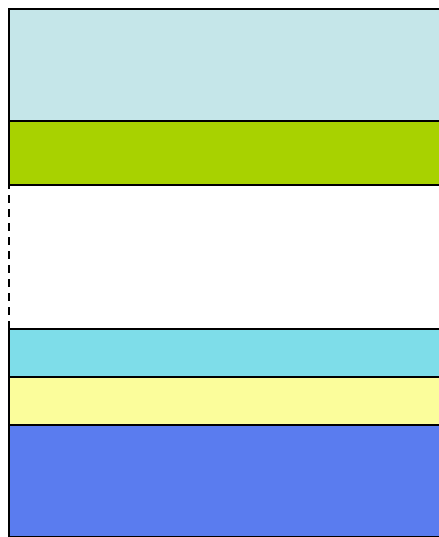


two central problems

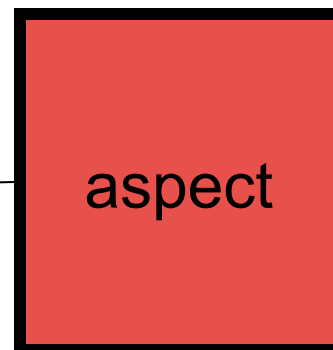
AOP tries to solve

code tangling:

one module
many concerns

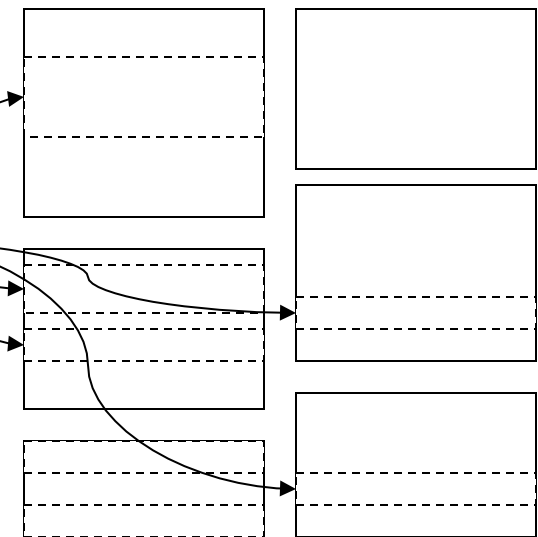


example:
logging



code scattering:

one concern
many modules



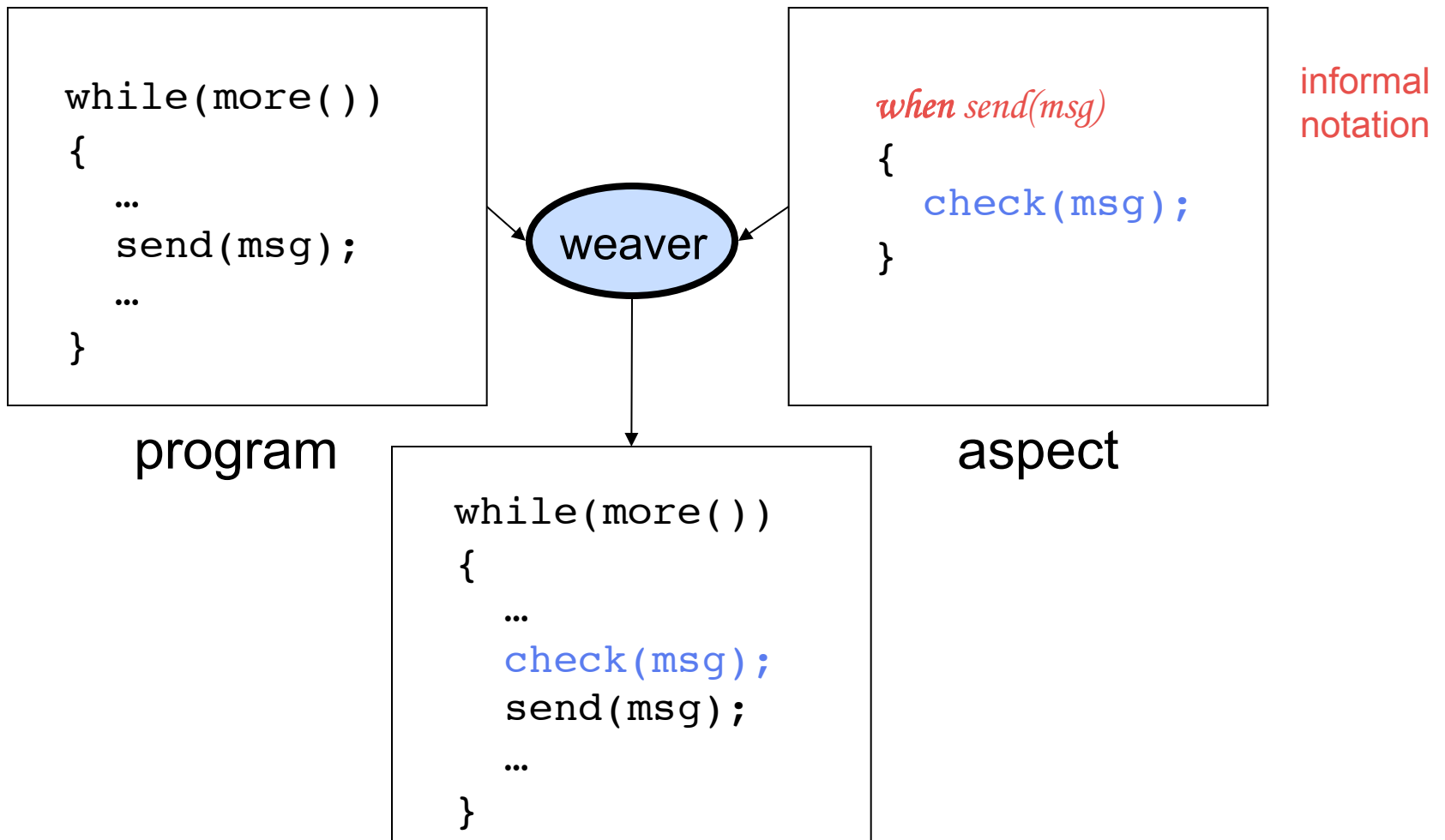
examples of crosscutting code

- logging (tracking program behavior)
- verification (checking program behavior)
- policy enforcement (correcting behavior)
- security management (preventing attacks)
- profiling (exploring where a program spends its time)
- memory management
- visualization of program executions
- ...

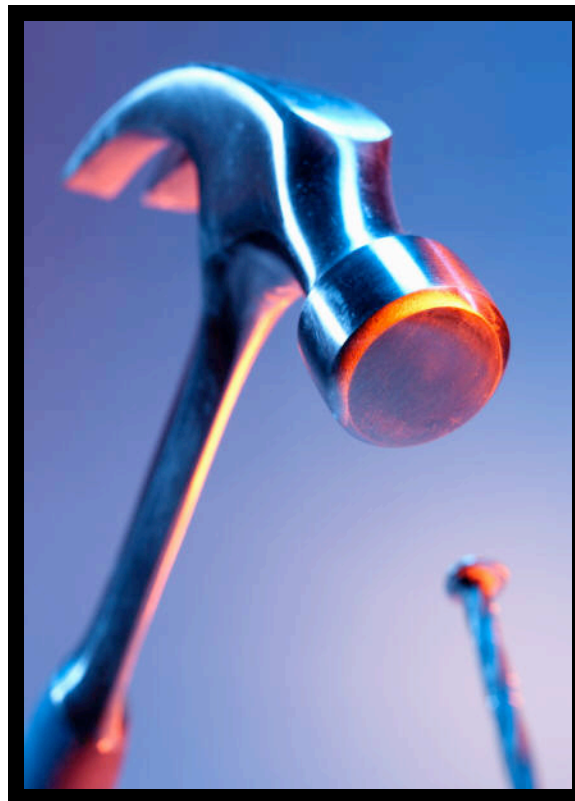
the problem

- the flow of the core logic gets obscured, harder to follow, the core logic is **tangled** with the new code.
- the new code code gets **scattered** throughout the code base
 - lots of typing
 - big picture (in one place) is missing
 - difficult to find what is new code and how it works
 - difficult to change new code
 - increases probability of consistency errors

very simplified view of AOP



that's it



except for
notation,
all the
details,
usage,
...

basic mechanisms

- join points
 - points in a Java program
- three main additions to Java
 - pointcut
 - picks out join points and values at those points
 - primitive and user-defined pointcuts
 - advice
 - additional action to take at join points matching a pointcut
 - aspect
 - a modular unit of crosscutting behavior
 - normal Java declarations
 - pointcut definitions
 - advice

- inter-type declarations
add fields, methods to classes

terminology as equations

Program:

Joinpoint = well-defined point in the program

AspectJ:

Pointcut = **Joinpoint**-set

Advice = Kind \times **Pointcut** \times Code
where Kind = {before, after, around}

Aspect = **Advice**-list

example

```
class Power {  
    int balance;  
  
    void deposit(int amount) {  
        balance = balance + amount;  
    }  
  
    boolean withdraw(int amount) {  
        if (balance - amount > 0) {  
            balance = balance - amount;  
            return true;  
        } else return false;  
    }  
}
```

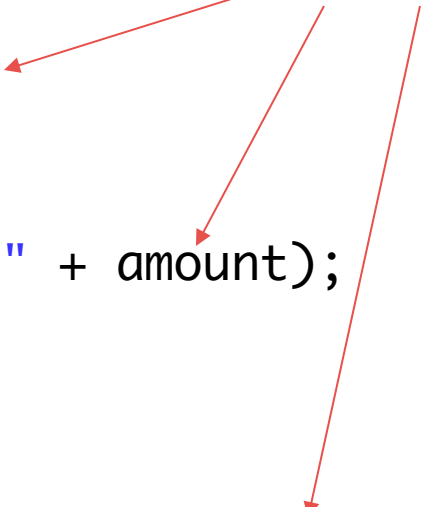
logging class

```
class Logger {  
    private PrintStream stream;  
  
    Logger() {  
        ... create stream  
    }  
  
    void log(String message) {  
        stream.println(message);  
    }  
}
```

logging the traditional way

```
class Power {  
    int balance;  
    Logger logger = new Logger();  
  
    void deposit(int amount) {  
        logger.log("deposit amount: " + amount);  
        balance = balance + amount;  
    }  
  
    boolean withdraw(int amount) {  
        logger.log("withdraw amount: " + amount);  
        if (balance - amount >= 0) {  
            balance = balance - amount;  
            return true;  
        } else return false;  
    }  
}
```

logging

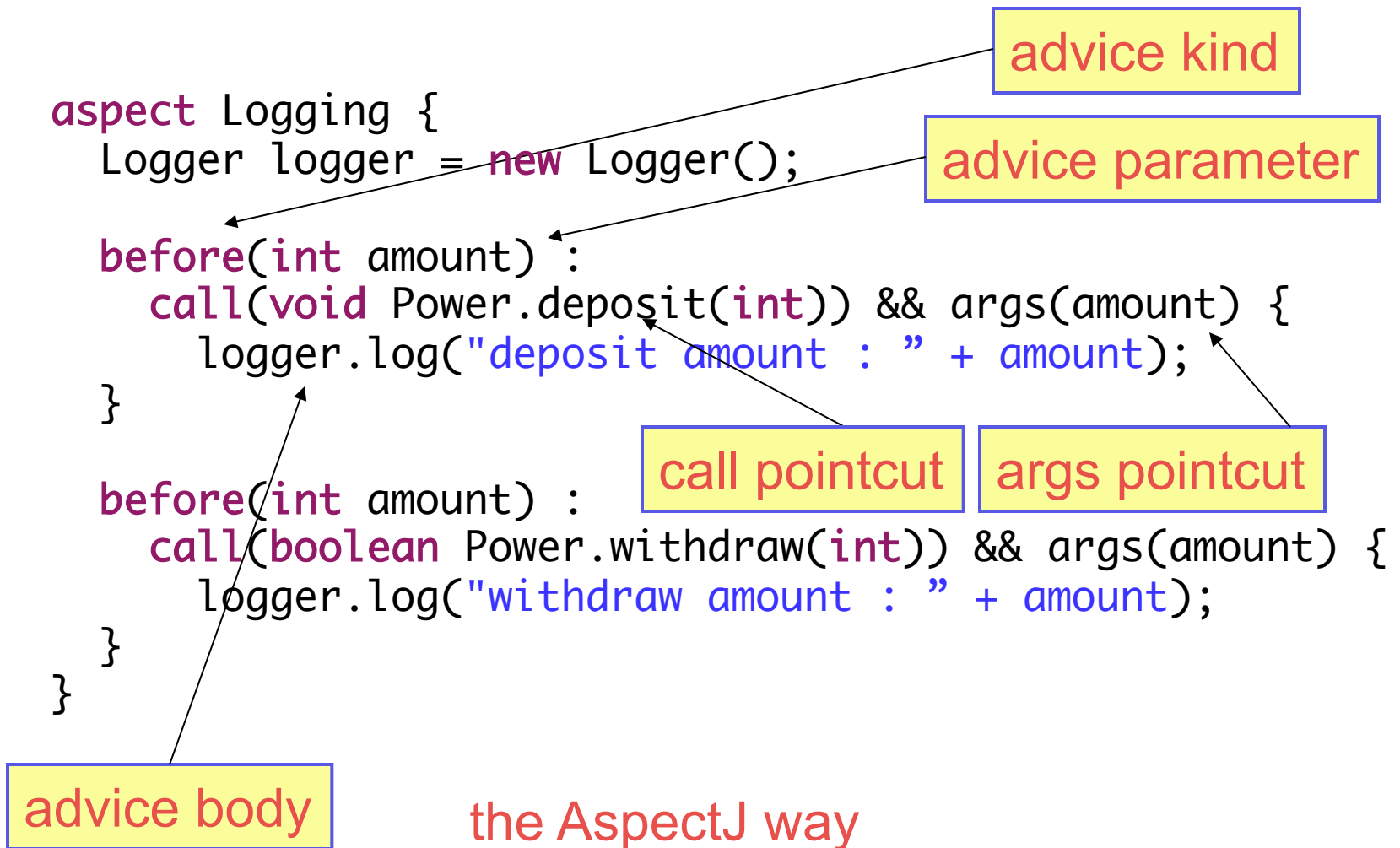


logging the AOP way

```
aspect Logging {  
    Logger logger = new Logger();  
  
    when deposit(amount) {  
        logger.log("deposit amount : " + amount);  
    }  
  
    when withdraw(amount) {  
        logger.log("withdraw amount : " + amount);  
    }  
}
```

that's not quite how it is written though

logging the AOP way



primitive pointcuts

a pointcut is a predicate on join points that:

- can match or not match any given join point and
- optionally, can pull out some of the values at that join point

Example:

```
call(void Power.deposit(int))
```

matches any join point that is a:
call of a method with this signature


explaining parameters...

of advice

- variable is bound by advice declaration
 - pointcut supplies value for variable
 - value is available in advice body

advice parameter

typed variable in place
of type name

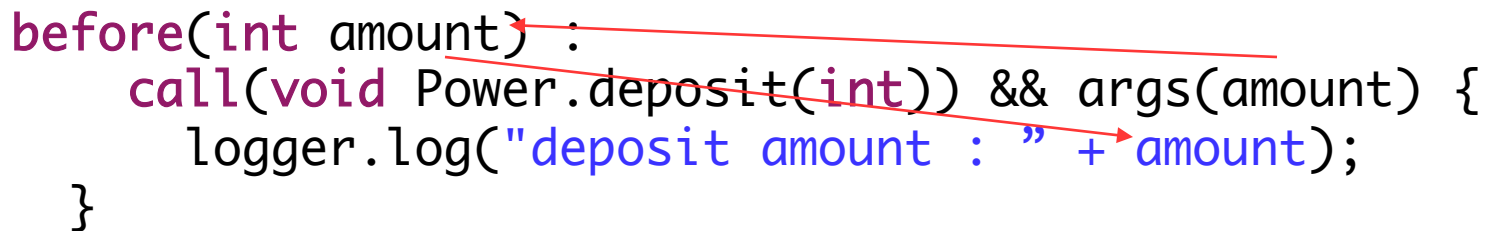


```
before(int amount) :  
    call(void Power.deposit(int)) && args(amount) {  
        logger.log("deposit amount : " + amount);  
    }
```

parameter data flow

- value is 'pulled'
 - right to left across ':' from pointcuts to advice
 - and then to advice body

```
before(int amount) :  
    call(void Power.deposit(int)) && args(amount) {  
        logger.log("deposit amount : " + amount);  
    }
```

A red arrow originates from the 'amount' parameter in the pointcut 'before(int amount) :' and points to the 'amount' parameter in the advice body 'logger.log("deposit amount : " + amount);', illustrating the 'pulling' of the value from the pointcut to the advice body.

pointcut naming and patterns

named pointcut

```
aspect Balance {
```

```
    pointcut powerChange(Power power) :  
        (call(* deposit(..)) || call(* withdraw(..)))  
        && target(power);
```

pointcut patterns

```
    after(Power power) : powerChange(power) {  
        System.out.println("balance = " + power.balance);  
    }
```

```
}
```

"after" advice

target pointcut

privileged aspects

can access private fields and methods

→ **privileged aspect** Balance {

```
pointcut powerChange(Power power) :  
    (call(* deposit(..)) || call(* withdraw(..)))  
    && target(power);
```

```
after(Power power) : powerChange(power) {  
    System.out.println("balance = " + power.balance);  
}  
}
```

suppose power.balance is a private variable. Then the aspect must be **privileged**.

args, this and target pointcuts

```
before(Rover rover, Power power, int amount) :  
    call(void Power.deposit(int))  
    && args(amount) && this(rover) && target(power) {...}
```

Object R

Object P

```
class Rover {  
    ...  
    void execute(...) {  
        ...  
        power.deposit(500);  
        ...  
    }  
    ...  
}
```

```
class Power {  
    ...  
    void deposit(int amount){  
        ...  
    }  
    ...  
}
```

target pointcut

target(*TypeName* | *VariableName*)

does two things:

- predicate on join points - any join point at which target object is an instance of *TypeName* or of the same type as *VariableName*.
“any join point “ can be:
 - method call join points
 - field get & set join points
 - ...
- exposes target if argument is a variable name

target(Power) :

- matches when target object is of type Power

Power is a type

target(power) :

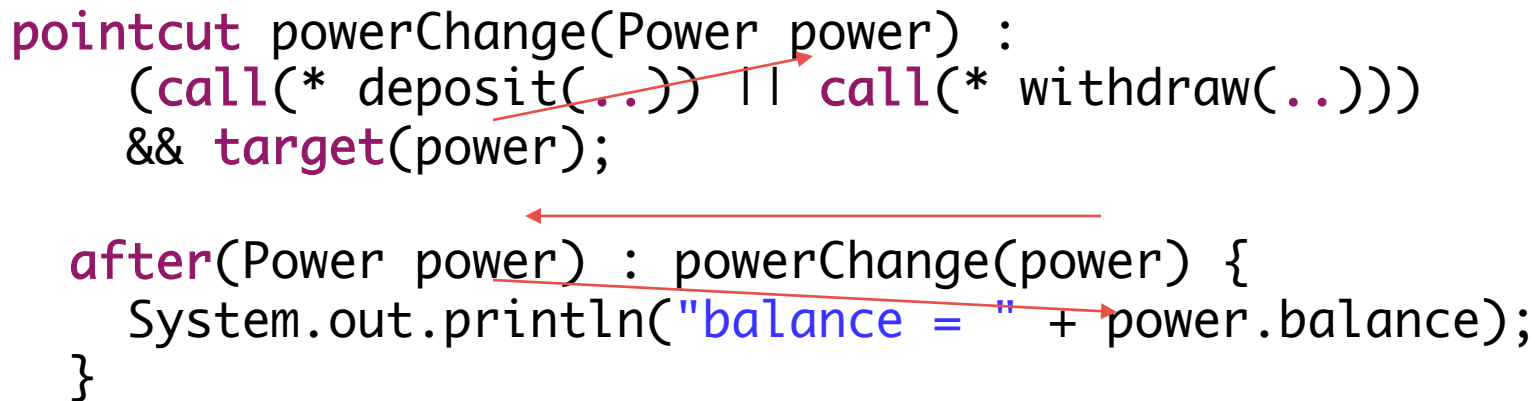
- ditto, since power is of type Power
- in addition it binds the target object to power

power is a variable

parameter data flow again

- value is 'pulled'
 - right to left across ':' from pointcuts to user-defined pointcuts
 - from pointcuts to advice
 - and then to advice body

```
pointcut powerChange(Power power) :  
    (call(* deposit(..)) || call(* withdraw(..)))  
    && target(power);  
  
after(Power power) : powerChange(power) {  
    System.out.println("balance = " + power.balance);  
}
```



contract checking

- pre-conditions
 - check that parameter is valid
- post-conditions
 - check that result is correct
- policy enforcement
 - check, and correct if check fails
- invariants
 - check that some condition on the state
“always” holds

pre-condition

fictive notation

```
boolean withdraw(int amount) {  
    pre balance - amount > 50;  
    ...  
    // implementation:  
    ...  
}
```

using **before** advice

pre-condition

```
aspect WithDrawPreCond {  
    final int MIN_BALANCE = 50;  
  
    before(Power power, int amount) :  
        call(boolean Power.withdraw(int)) &&  
        target(power) && args(amount)  
    {  
        assert power.balance - amount > MIN_BALANCE :  
            "withdrawal too big: " + amount;  
    }  
}
```

post condition

fictive notation

```
boolean withdraw(int amount) {  
    pre balance - amount > 50;  
    post result == (old(balance) - amount) >= 0  
        &&  
        balance ==  
            (result ? old(balance) - amount  
                : old(balance))  
    ...  
    // implementation:  
    ...  
}
```

using **before** & **after** advice

post-condition

```
public aspect WithdrawPostCond {  
    int old_balance;  
  
    before(Power power) :  
        call(boolean Power.withdraw(int)) && target(power)  
    {  
        old_balance = power.balance;  
    }  
  
    after(Power power, int amount) returning(boolean changed) :  
        call(boolean Power.withdraw(int)) &&  
        target(power) && args(amount)  
    {  
        assert changed == (old_balance - amount) >= 0 &&  
            power.balance ==  
                (changed ? old_balance-amount : old_balance);  
    }  
}
```

post-condition

around advice

```
aspect WithDrawPostCondAround {  
    int old_balance;  
  
    boolean around(Power power, int amount) :  
        call(boolean Power.withdraw(int)) &&  
        target(power) && args(amount)  
    {  
        old_balance = power.balance;  
        boolean changed = proceed(power, amount);  
        assert changed == (old_balance - amount) >= 0;  
        assert power.balance ==  
            (changed ? old_balance - amount : old_balance);  
        return changed;  
    }  
}
```

proceed statement

the proceed “method”

for each around advice with the signature:

T around(*T1* arg1, *T2* arg2, ...)

there is a special method with the signature:

T proceed(*T1*, *T2*, ...)

calling this method means:

“run what would have run if this around advice had not been defined”

policy enforcement

fictive notation

```
boolean withdraw(int amount) {  
    pre balance - amount > 50 {  
        System.out.println("withdrawal rejected");  
        return false;  
    }  
    ...  
    // implementation:  
    ...  
}
```

policy enforcement

around advice

```
aspect WithdrawCorrect {  
    final int MIN_BALANCE = 50;  
  
    boolean around(Power power, int amount) :  
        call(boolean Power.withdraw(int)) &&  
        target(power) && args(amount)  
    {  
        if(power.balance - amount >= MIN_BALANCE)  
            return proceed(power, amount);  
        else {  
            System.out.println("withdrawal rejected");  
            return false;  
        }  
    }  
}
```

proceed statement

invariant checking

fictive notation

```
class Power {  
    int balance;  
  
    invariant balance >= 500;  
  
    void deposit(int amount) {  
        ...  
    }  
  
    boolean withdraw(int amount) {  
        ...  
    }  
}
```

invariant checking

```
aspect Invariant {  
    boolean invariant(int balance) {  
        return balance >= 500;  
    }  
  
    pointcut write(int balance) :  
        set(int Power.balance) && args(balance);  
  
    before(int balance) : write(balance) { //every update  
        if (!invariant(balance))  
            System.out.println("invariant violated");  
    }  
}  
  
    after(Power power) :  
        execution(* Power.*(..)) && target(power)  
    {  
        if (!invariant(power.balance))  
            System.out.println("invariant violated");  
    } // at method boundaries
```

examples of patterns

Type names:

Command
*Command
java.*.Date
Java..*
Javax..*Model+

Combined Types:

!Vector
Vector || HashTable
java.util.RandomAccess+
&& java.util.List+

Method Signatures:

public void Power.set*(*)
boolean Power.withdraw(int)
bo* Po*.wi*w(i*)
!static * *.*(...)
rover..command.Command+.check(int,...)

reflexive information available at **all** joinpoints

- **thisJoinPoint**

- getArgs() : Object[]
- getTarget() : Object
- getThis() : Object
- getStaticPart() : JoinPointStaticPart



- **thisJoinPointStaticPart**

- getKind() : String
- getSignature() : Signature
- getSourceLocation() : SourceLocation

logging exceptions using thisJoinPoint

```
aspect LogExceptions {  
    Logger logger = new Logger();  
  
    after() throwing (Error e): call(* *(..)) {  
        logger.log("exception thrown " +  
            thisJoinPoint + ":" + e);  
    }  
}
```

logged information in the case of an assertion error in call of withdraw:

```
...  
exception thrown call(boolean core.Power.withdraw(int)):java.lang.AssertionError  
...  
                thisJoinPoint
```

checking object creation

```
class CmdFactory {  
    static Command mkTurnCommand(int budget, int degrees) {  
        return new TurnCommand(budget, degrees);  
    }  
    ...  
}
```

want to ensure that any creation of commands goes through the factory methods mk...

```
aspect FactoryCheck {  
    pointcut illegalNewCommand():  
        call(Command+.new(..)) &&  
        !withincode(* CmdFactory.mk*(..));  
  
    before(): illegalNewCommand() {  
        throw new Error("Use factory method instead.");  
    }  
}
```

checking object creation

```
class CmdFactory {  
    static Command mkTurnCommand(int budget, int degrees) {  
        return new TurnCommand(budget, degrees);  
    }  
    ...  
}
```

want to ensure that any creation of commands goes through the factory methods mk...

```
aspect FactoryCheckStatic {  
    pointcut illegalNewCommand():  
        call(Command+.new(..)) &&  
        !withincode(* CmdFactory.mk*(..));  
  
    declare error : illegalNewCommand() :  
        "Use factory method instead."  
}
```

must be a "static pointcut"

static
declare

causes check to be
performed at compile
time

inter-type declarations

- inside an aspect:
adding declarations to a class **C**

```
aspect A {  
  int counter = 0;  
  void count() {counter++;}  
  ...  
}
```



```
class C {  
  ...  
  ...  
}
```

inter-type declarations

- one must indicate what class:

```
aspect A {  
    int C.counter = 0;  
    void C.count() {counter++;}  
    ...  
}
```




```
class C {  
    ...  
    ...  
}
```

inserting fields and methods

verify that a command is executed no more than once!

requires a counter per Command object.

field and method inserted in Command object but accessible only to aspect.



```
aspect ExecuteOnlyOnce {  
    private int Command.counter = 0;  
    private void Command.count() {counter++;}  
  
    before(Command cmd) :  
        call(void Command+.execute()) && target(cmd)  
    {  
        assert cmd.counter == 0 : "command executed again";  
        cmd.count();  
    }  
}
```

same property

verify that a command is executed no more than once!

eliminating: **private**, **+**, count method, assert message


```
aspect ExecuteOnlyOnce {  
    int Command.counter = 0;  
  
    before(Command cmd) :  
        call(void Command.execute()) && target(cmd)  
    {  
        assert cmd.counter++ == 0;  
    }  
}
```

it does not get much shorter than this

aspect association

- instances of aspects:
 - one per virtual machine (the default)
 - one per object (*perthis*, *pertarget*)
 - one per control-flow (*percfow*, *percfowbelow*)

```
aspect <AspectName> <association>(<pointcut>){  
    pointcut ... : ...;  
    ...  
}  
  
<association> ::=  
    perthis | pertarget | percfow | percfowbelow
```



aspect association

- `perthis(pc)`:
 - when a pointcut satisfying `pc` is reached, and `this(x)` holds, and `x` does not already have an associated aspect instance of this type, a new instance is created for `x` (to track `x`)
- `pertarget(pc)`:
 - similar, except we use `target(x)`
- `percflow(pc)`:
 - when a pointcut satisfying `pc` is reached, a new instance is created, which lasts as long as the control flow under this `pc` does

same property


verify that a command is executed no more than once!

this time using object association : one aspect per Command
target of the execute command.

```
aspect ExecuteOnlyOnce2 pertarget(execute()){  
    int counter = 0;  
  
    pointcut execute() : call(void Command.execute());  
  
    before() : execute() {  
        assert counter++ == 0;  
    }  
}
```

Tracing aspect

all pointcuts
except within
Tracing aspect



```
aspect Tracing {  
    private int callDepth = -1;  
  
    pointcut tracePoint() : !within(Tracing);  
  
    before() : tracePoint() {  
        callDepth++; print("Before", thisJoinPoint);  
    }  
  
    after() : tracePoint() {  
        print("After", thisJoinPoint); callDepth--;  
    }  
  
    private void print(String prefix, Object message) {  
        for(int i = 0, spaces = callDepth * 2; i < spaces; i++)  
            System.out.print(" ");  
        System.out.println(prefix + ": " + message);  
    }  
}
```

```
<terminated> Main (2) [Java Application] /System/Library/Frameworks/JavaVM.framework/Versions/1.5.0/Home/bin/java (Mar 25, 2008 10:15:52 AM)
Before: execution(core.ExecuteOnlyOnce())
After: execution(core.ExecuteOnlyOnce())
After: initialization(core.ExecuteOnlyOnce())
After: staticinitialization(core.ExecuteOnlyOnce.<clinit>)
Before: set(int core.Command.counter)
After: set(int core.Command.counter)
Before: execution(core.Command(int))
Before: set(int core.Command.budget)
After: set(int core.Command.budget)
After: execution(core.Command(int))
After: initialization(core.Command(int))
Before: initialization(core.PictureCommand(int))
Before: execution(core.PictureCommand(int))
After: execution(core.PictureCommand(int))
After: initialization(core.PictureCommand(int))
After: call(core.PictureCommand(int))
Before: call(core.TurnCommand(int, int))
Before: staticinitialization(core.TurnCommand.<clinit>)
After: staticinitialization(core.TurnCommand.<clinit>)
Before: preinitialization(core.TurnCommand(int, int))
After: preinitialization(core.TurnCommand(int, int))
Before: preinitialization(core.Command(int))
After: preinitialization(core.Command(int))
Before: initialization(core.Command(int))
Before: set(int core.Command.counter)
After: set(int core.Command.counter)
Before: execution(core.Command(int))
Before: set(int core.Command.budget)
After: set(int core.Command.budget)
After: execution(core.Command(int))
After: initialization(core.Command(int))
Before: initialization(core.TurnCommand(int, int))
Before: execution(core.TurnCommand(int, int))
Before: set(int core.TurnCommand.degrees)
After: set(int core.TurnCommand.degrees)
```

output

abstract pointcuts

- what if we want to trace specific events? do we edit the Tracing aspect? no, we can define the pointcut as abstract
- a pointcut can be defined as abstract without a “right-hand” side:

```
abstract pointcut something(T x);
```

- advices can be defined on the abstract pointcut
- specialization of aspect can later define the pointcut
- this resembles parameterization with pointcuts
- similar to the way methods can be defined abstract and later defined in sub-classes

abstract Tracing aspect

```
abstract aspect AbstractTracing {  
    private int callDepth = -1;  
  
    abstract pointcut tracePoint();  
  
    before() : tracePoint() {  
        callDepth++; print("Before", thisJoinPoint);  
    }  
  
    after() : tracePoint() {  
        print("After", thisJoinPoint); callDepth--;  
    }  
  
    private void print(String prefix, Object message) {  
        for(int i = 0, spaces = callDepth * 2; i < spaces; i++)  
            System.out.print(" ");  
        System.out.println(prefix + ": " + message);  
    }  
}
```

aspect and pointcut are now abstract, the rest is the same!

concrete tracing aspect

we just define
the pointcut

```
aspect ConcreteTracing extends AbstractTracing {  
    pointcut tracePoint() :  
        call(* Power.*(..)) || withincode(* Command+.*(..));  
}
```

It's a bit like function application:

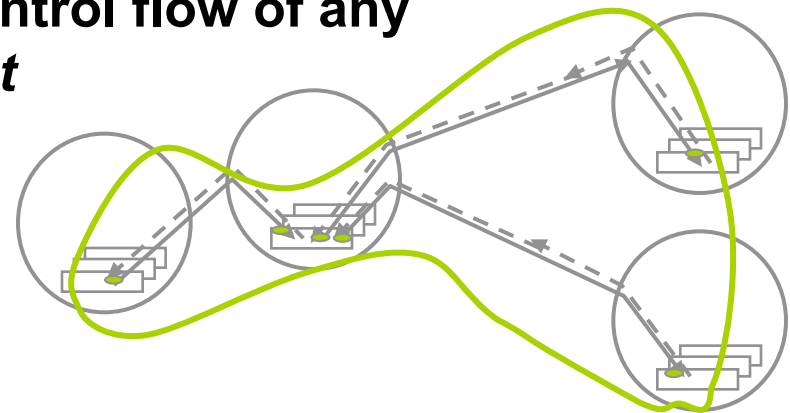
```
aspect ConcreteTracing =  
    AbstractTracing(  
        call(* Power.*(..)) || withincode(* Command+.*(..))  
    )
```

not AspectJ syntax

control flow pointcuts

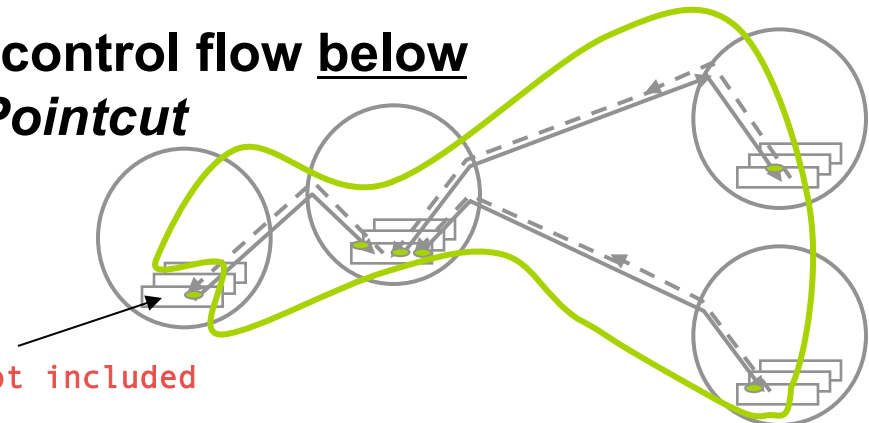
cflow(*Pointcut*)

all join points in the dynamic control flow of any join point picked out by *Pointcut*



cflowbelow(*Pointcut*)

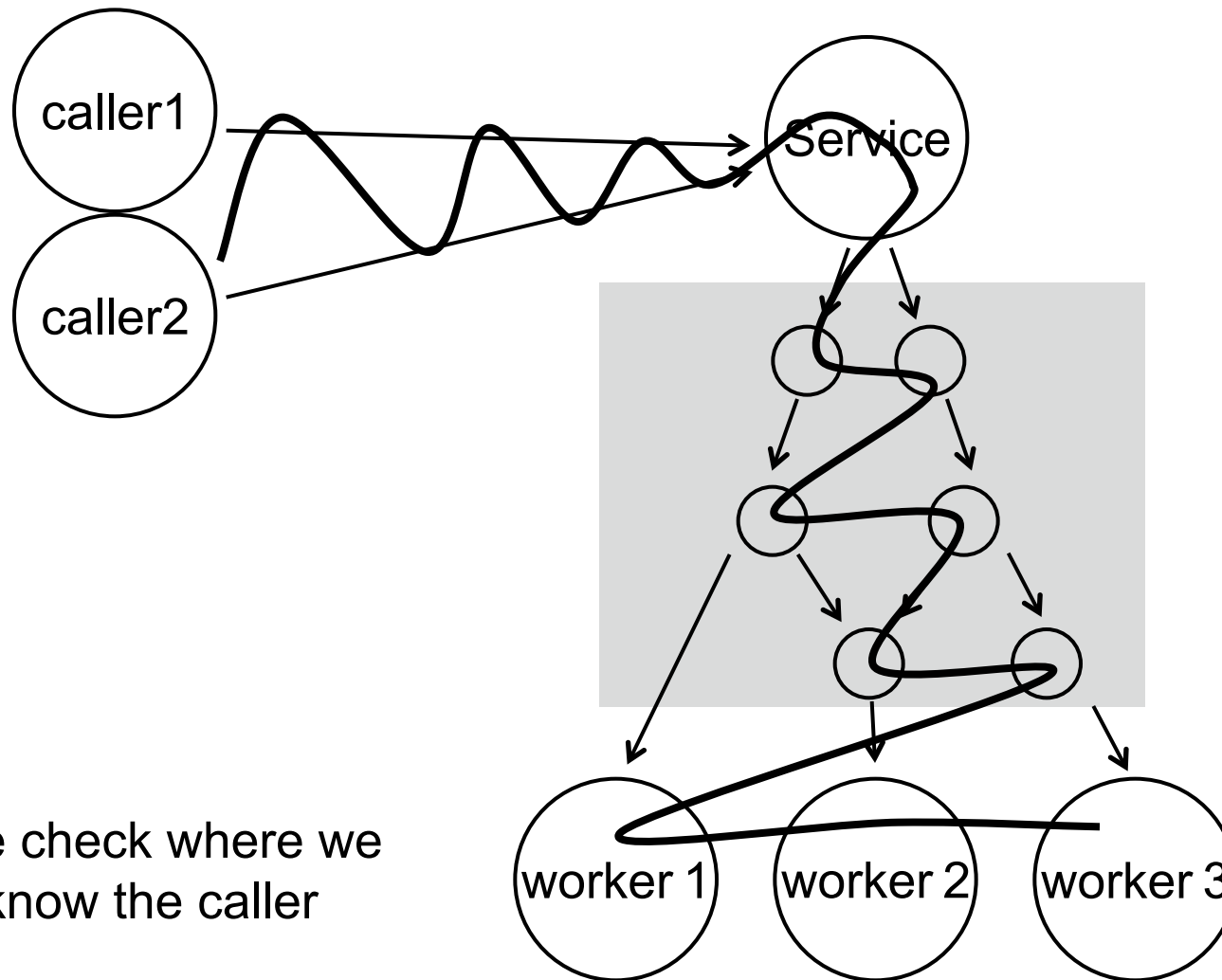
all join points in the dynamic control flow below any join point picked out by *Pointcut*



top pointcut not included

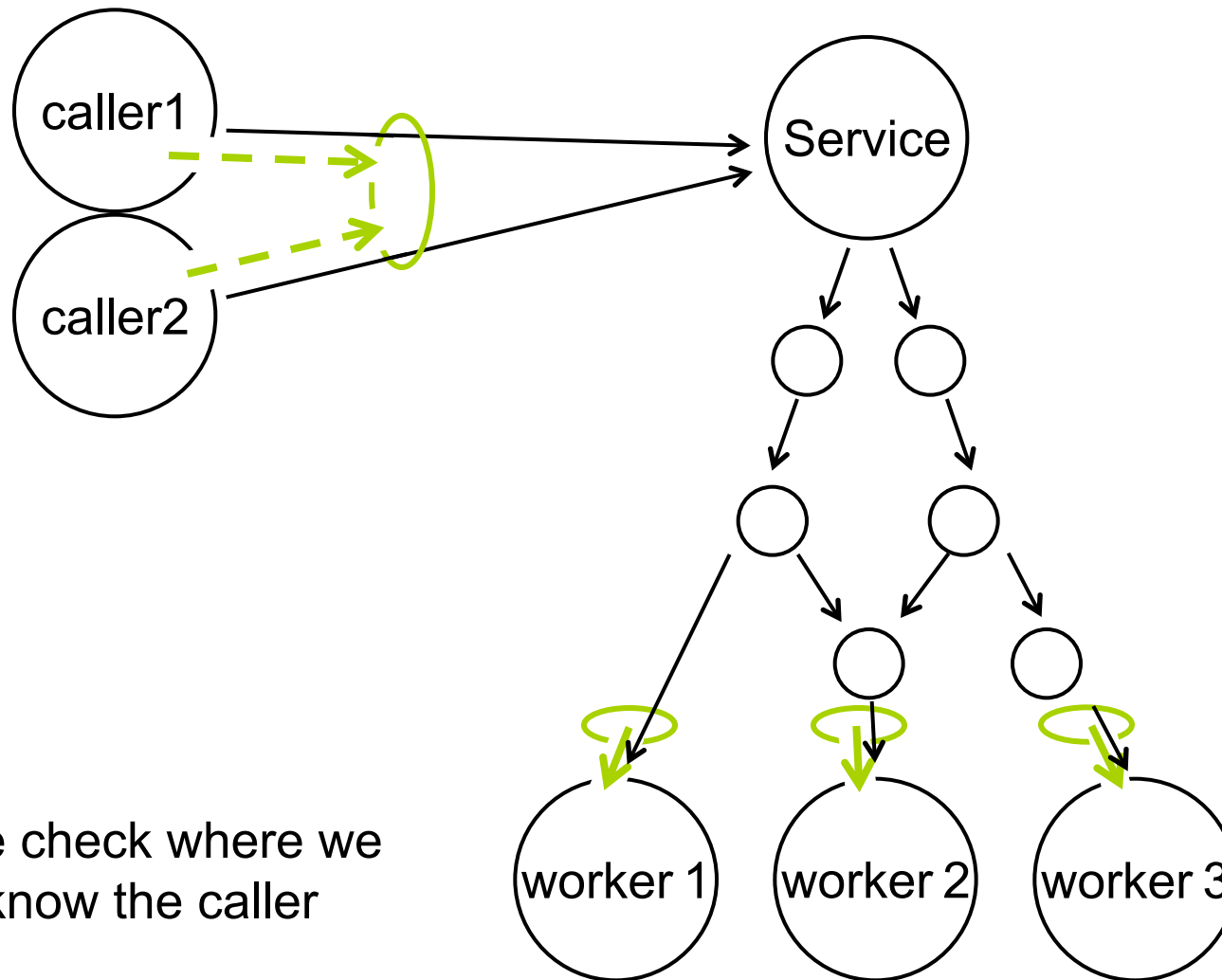
context-passing aspects

example



introduce check where we
need to know the caller

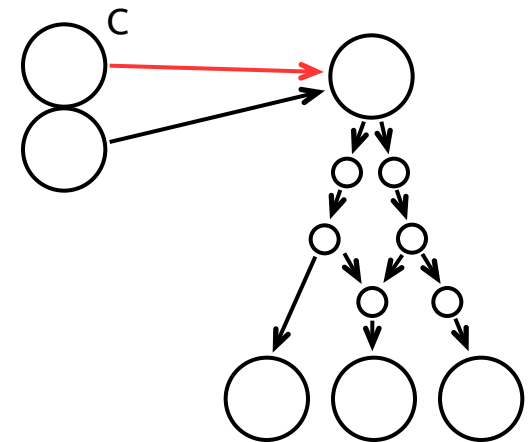
example



introduce check where we
need to know the caller

context-passing aspects

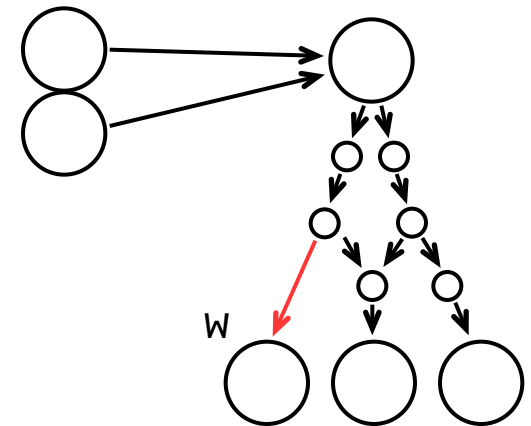
```
pointcut invocation(Caller c):  
    this(c) && call(void Service.doService(String));
```



context-passing aspects

```
pointcut invocation(Caller c):  
    this(c) && call(void Service.doService(String));
```

```
pointcut workPoint(Worker w):  
    target(w) && call(void Worker.doTask(Task));
```

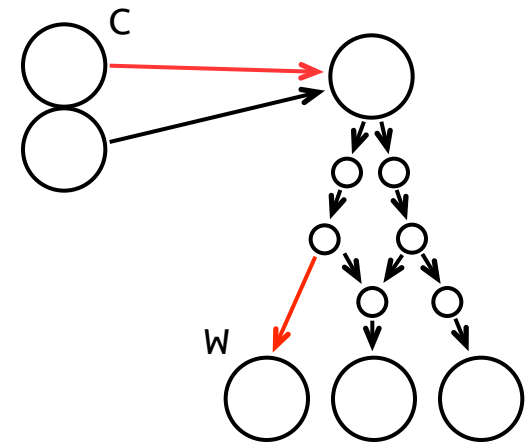


context-passing aspects

```
pointcut invocation(Caller c):  
    this(c) && call(void Service.doService(String));
```

```
pointcut workPoint(Worker w):  
    target(w) && call(void Worker.doTask(Task));
```

```
pointcut calledWork(Caller c, Worker w):  
    cflow(invocation(c)) && workPoint(w);
```



context-passing aspects

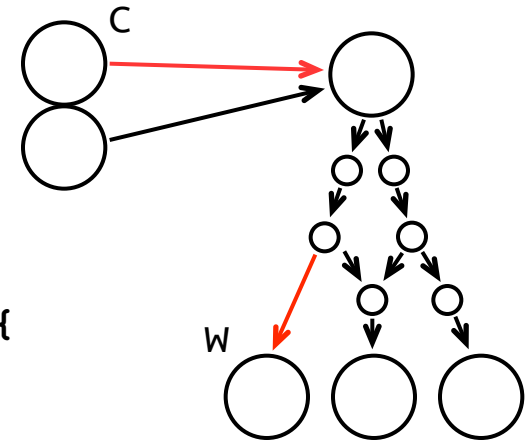
```
abstract aspect CapabilityChecking {
```

```
    pointcut invocation(Caller c):  
        this(c) && call(void Service.doService(String));
```

```
    pointcut workPoint(Worker w):  
        target(w) && call(void Worker.doTask(Task));
```

```
    pointcut calledWork(Caller c, Worker w):  
        cflow(invocation(c)) && workPoint(w);
```

```
    before (Caller c, Worker w): calledWork(c, w) {  
        verifyCalledWork(c, w);  
    }  
}
```



advice precedence

what happens if two pieces of advice apply to the same join point?

```
aspect Policy {  
    pointcut scope() : !cflow(adviceexecution());  
    ...  
    before(): call(* *.*(..)) && scope() {  
        if (!isAllowed(thisJoinPoint))  
            error("invalid ");  
    }  
  
    declare precedence: Policy, *;  
}
```

order undefined unless:

- in same aspect,
- in sub-aspect, or
- using declare precedence...

```
aspect LogIt {  
    pointcut scope() : !cflow(adviceexecution());  
  
    before(): call(* *.*(..)) && scope() {  
        System.out.println("Entering " + thisJoinPoint);  
    }  
}
```

advice precedence rules

assume that aspect **L** has
lower priority than aspect **H** ($L < H$)
and consider a particular joinpoint

- **H** executes its before advice **before**
L's before advice
- **H** executes its after advice **after**
L's after advice
- **H's** around advice **encloses**
L's around advice

beginner mistake

not controlling circularity of advice

pointcuts sometimes match more than expected

```
aspect A {  
    before(): call(String toString()) {  
        System.err.println(thisJoinPoint);  
    }  
}
```

use within, cflow, adviceexecution() to control

```
aspect A {  
    before(): call(String toString()) && !within(A) {  
        System.err.println(thisJoinPoint);  
    }  
}
```

summary

pointcuts

primitive

call
execution
handler
get set
initialization
this target args
within withincode
cflow cflowbelow

user-defined

pointcut

advice

before
after
around

inter-type decls

Type.field
Type.method()

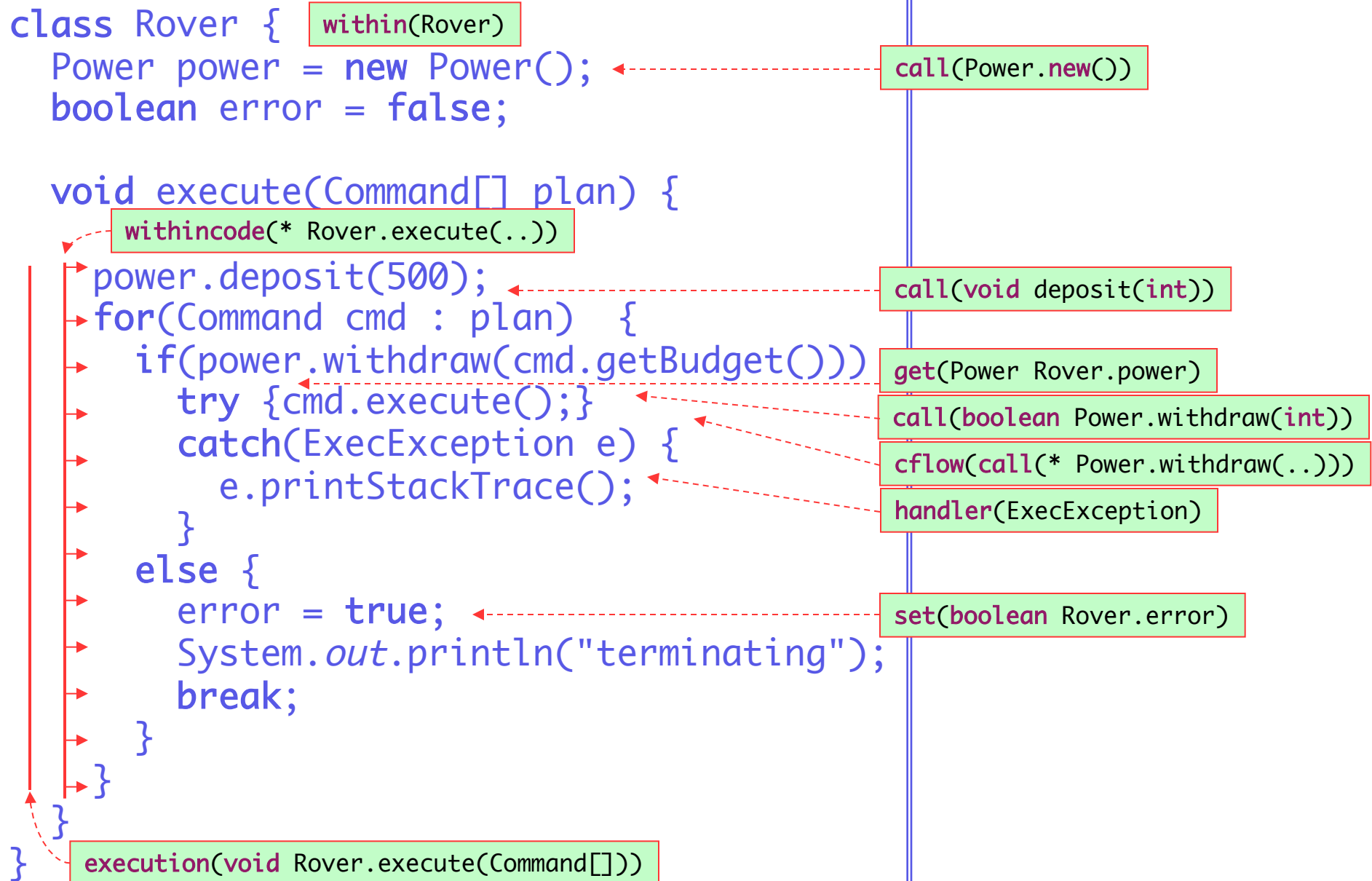
declare

error
parents
precedence

reflection

thisJoinPoint
thisJoinPointStaticPart

pointcut overview (see aspect next two slides)



an aspect that gets “around”

```
public aspect Monitor {
    static boolean tracingOn = true;

    pointcut scope() : if(tracingOn) && !cflow(adviceexecution());

    pointcut handlethrow(ExecException e) : handler(ExecException) && args(e);
    before(ExecException e) : handlethrow(e) && scope() {
        print("*** bad luck: " + e);
    }

    after() returning (Power power) : call(Power.new()) && scope() {
        print("power object created " + power);
    }

    before(int amount) : call(void deposit(int)) && args(amount) && scope() {
        print("depositing: " + amount);
    }

    after(int amount) returning (boolean success):
        call(boolean Power.withdraw(int)) && args(amount) && scope() {
            print("withdrawing " + amount + ":" + success);
        }
}
```

```

void around(Command[] plan) :
    execution(void Rover.execute(Command[])) && args(plan) && scope() {
        if (!validatePlan(plan))
            proceed(correctPlan(plan));
        else
            proceed(plan);
    }

after() returning(Power power):
    get(Power Rover.power) && within(Rover) && scope() {
        print("reading power " + power);
    }

before(boolean value) :
    set(boolean Rover.error) && args(value) && if(value)
    && withcode(* Rover.execute(..)) && scope() {
        print("error flag being set to " + value);
    }

before() : call(* *.*(..)) && cflow(call(* Power.withdraw(..))) && scope() {
    print("function call " + thisJoinPointStaticPart.getSignature());
}

before(Rover rover, Command command) :
    call(* Command.execute()) &&
    this(rover) && target(command) && scope() {
        print("Rover " + rover + " executing command " + command);
    }
}

```

... continued

abstract syntax for AspectJ

- contains most elements of language
- look at quick guide
- look at examples

AspectDecl ::=
[privileged] [Modifiers] aspect Id
[extends Type] [implements TypeList]
[PerClause]
{ BodyDecl* }

PerClause ::=
pertarget (Pointcut) | perthis (Pointcut)
| perflow (Pointcut) | perflowbelow (Pointcut) | issingleton ()

BodyDecl ::=
JavaBodyDecl
| IntertypeDecl
| PointcutDecl
| AdviceDecl

InterTypeDecl ::=

- [Modifiers] Type Type . Id (Formals) [throws TypeList] { Body }
- | [Modifiers] Type . new (Formals) [throws TypeList] { Body }
- | [Modifiers] Type Type . Id [= Expression] ;
- | declare warning : Pointcut : String ;
- | declare error : Pointcut : String ;
- | declare precedence : TypePatList ;

PointcutDecl ::=

abstract [Modifiers] **pointcut** Id (Formals) ;
 | [Modifiers] **pointcut** Id (Formals) : Pointcut ;

AdviceDecl ::=

 AdviceSpec [**throws** TypeList] : Pointcut { Body }

AdviceSpec ::=

before (Formals)
 after (Formals)
 after (Formals) returning [(Formal)]
 after (Formals) throwing [(Formal)]
 Type **around** (Formals)

Pointcut ::=

`call(MethodPat) | call(ConstructorPat)`
`| execution(MethodPat) | execution(ConstructorPat)`
`| initialization(ConstructorPat) | preinitialization(ConstructorPat)`
`| staticinitialization(TypePat)`
`| get(FieldPat) | set(FieldPat)`
`| handler(TypePat)`
`| adviceexecution()`
`| within(TypePat) | withincode(MethodPat) | withincode(ConstructorPat)`
`| cflow(Pointcut) | cflowbelow(Pointcut)`
`| if(Expression)`
`| this(Type | Var) | target(Type | Var) | args(Type | Var , ...)`

MethodPat ::=
[ModifiersPat] TypePat [TypePat .] IdPat (TypePat | .., ...)
[throws ThrowsPat]

ConstructorPat ::=
[ModifiersPat] [TypePat .] new (TypePat | .. , ...)
[throws ThrowsPat]

FieldPat ::= [ModifiersPat] TypePat [TypePat .] IdPat

TypePat ::=
IdPat [+] [[] ...]
| ! TypePat
| TypePat && TypePat
| TypePat || TypePat
| (TypePat)

IdPat ::=
Java id with `*'s mixed in

AspectJ syntax

special expressions and statements

Expression ::=

- thisJoinPoint
- | thisJoinPointStaticPart
- | thisEnclosingJoinPointStaticPart

StatementExpression ::=

proceed (Arguments)

end