



Applied Numerical Optimization

Prof. Alexander Mitsos, Ph.D.

Parametric optimization and uncertainty

Optimization Under Uncertainty

- Consider **uncertain** variables/parameters \mathbf{y} , e.g., :
 - \mathbf{y} is a model parameter that can at best be **estimated**,
 - \mathbf{y} results from a **stochastic process**, which is not yet resolved.

- \mathbf{y} can affect objective function and constraints

$$\begin{array}{ll} \min_{\mathbf{x}} & f(\mathbf{x}, \mathbf{y}) \\ \text{s. t.} & \mathbf{c}(\mathbf{x}, \mathbf{y}) \leq \mathbf{0} \end{array} \quad (\text{PAR})$$

- \mathbf{y} is not an optimization variable
 - if \mathbf{y} is fixed then (PAR) is normal NLP in \mathbf{x}
- Alternative notions to handle uncertainty:
 - **Parametric Optimization**
 - Stochastic Programming
 - Robust Optimization

Parametric Optimization

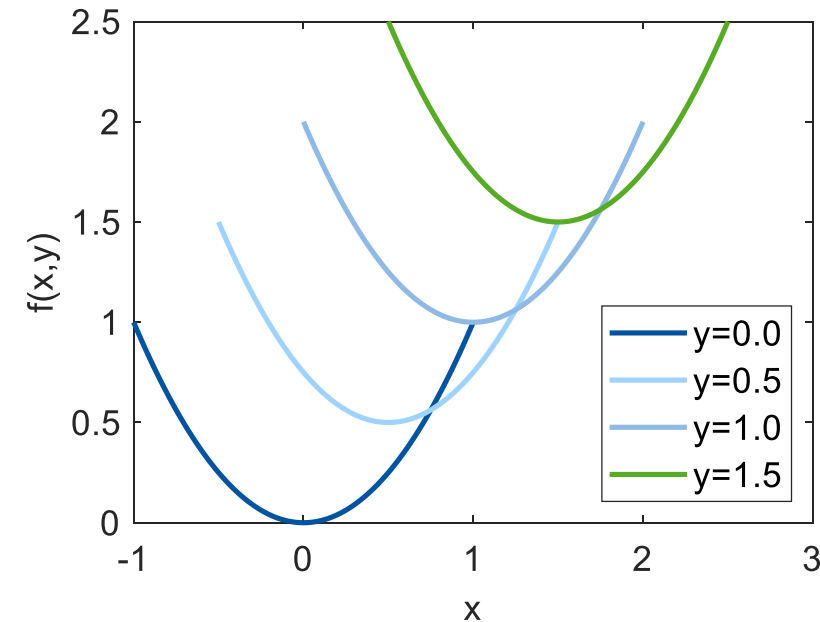
- Parametric optimization is finding solution as function of parameter

$$\begin{array}{ll} \min_x & f(x, y) \\ \text{s. t.} & c(x, y) \leq 0 \end{array} \quad (\text{PAR})$$

- The optimal solution point and optimal objective value depend on parameters: $x^*(y)$ and $f^*(y) = f(x^*(y), y)$

- Example:

- $\min_x (x - y)^2 + y$
- $x^*(y) = y$
- $f^*(y) = f(x^*(y), y) = (y - y)^2 + y = y$



$$\begin{aligned} f^*(\mathbf{y}) = \min_{\mathbf{x}} \quad & f(\mathbf{x}, \mathbf{y}) \\ \text{s. t.} \quad & \mathbf{c}(\mathbf{x}, \mathbf{y}) \leq \mathbf{0} \end{aligned} \quad (\text{PAR})$$

- Suppose that \mathbf{y} is uncertain parameter
- Parametric optimization is only useful under the following conditions:
 - The uncertainty is realized before the decision variables must be fixed
 - Once the uncertainty is realized, it is not practical to solve (PAR) for fixed \mathbf{y}These conditions hold for example [in online control](#) and [resource allocation](#).
- Solution approaches typically identify parameter regions for which the solution stays qualitatively same.
 - So called critical region (CR), e.g., no active set change

Example: Parametric Quadratic Optimization

Consider convex quadratic parametric optimization problem

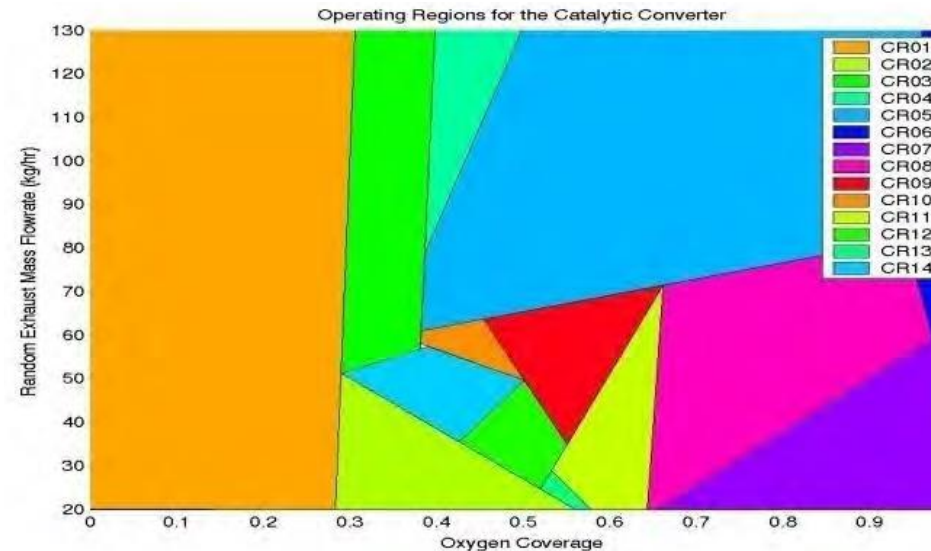
$$\begin{aligned} f^*(\mathbf{y}) = \min_{\mathbf{x} \in X} \quad & \frac{1}{2} \mathbf{x}^T \mathbf{G} \mathbf{x} + \mathbf{d}^T \mathbf{x} \\ \text{s. t.} \quad & \mathbf{A} \mathbf{x} \leq \mathbf{F} \mathbf{y} \end{aligned} \quad \text{with } \mathbf{y} \in Y$$

- Solve for $\mathbf{y} = \bar{\mathbf{y}}$
- Recall that 1st order KKT are necessary and sufficient
 - Under mild assumptions
- Recall sensitivity analysis via KKT conditions
 - Yields $\mathbf{x}^*(\cdot)$ as an affine function of \mathbf{y} .
 - $\mathbf{x}^*(\cdot)$ ceases to be valid once one of the inequalities in the KKT conditions is violated.

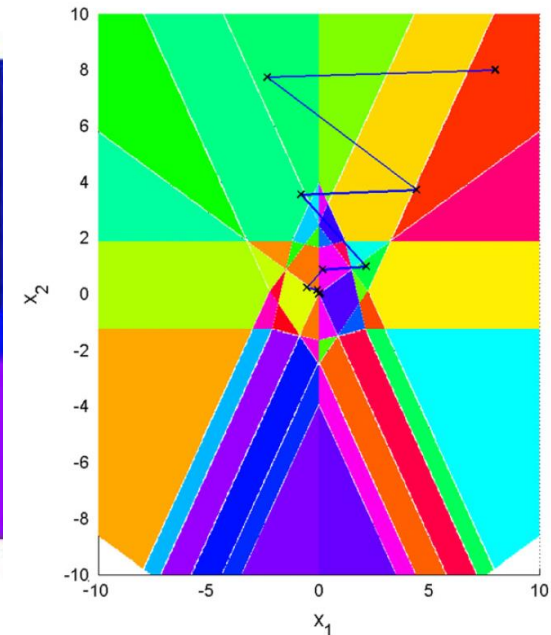
Example: Parametric Quadratic Optimization – Algorithm

1. Iterate $k = 1, 2, \dots$
2. Solve (QPAR) with $\mathbf{y} \in Y$ as free variable
3. Fix $\bar{\mathbf{y}}$ to the solution and solve (QPAR) for $\mathbf{y} = \bar{\mathbf{y}}$
4. Derive parametric solution for current critical region $CR^{(k)}$
5. Derive constraints to define $CR^{(k)}$
6. If Y is covered with critical regions, STOP
7. Go to step 1 with $CR^{(k)}$ removed from the feasible set

$$f^*(\mathbf{y}) = \min_{\mathbf{x} \in X} \frac{1}{2} \mathbf{x}^T \mathbf{G} \mathbf{x} + \mathbf{d}^T \mathbf{x} \quad (\text{QPAR})$$
$$\text{s. t. } \mathbf{A} \mathbf{x} \leq \mathbf{F} \mathbf{y}$$



Picture from: Pistikopoulos, OSE 2012



Picture from: Oberdieck, Pistikopoulos, Automatica 2015

Check Yourself

- What is parametric optimization and how does it relate to uncertainty?



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Introduction to optimization under uncertainty

Optimization Under Uncertainty

- Consider **uncertain** variables/parameters \mathbf{y} , e.g., :
 - \mathbf{y} is a model parameter that can at best be **estimated**,
 - \mathbf{y} results from a **stochastic process**, which is not yet resolved.

- \mathbf{y} can affect objective function and constraints

$$\begin{array}{ll} \min_{\mathbf{x}} & f(\mathbf{x}, \mathbf{y}) \\ \text{s. t.} & \mathbf{c}(\mathbf{x}, \mathbf{y}) \leq \mathbf{0} \end{array}$$

- \mathbf{y} is not an optimization variable
 - if \mathbf{y} is fixed then we have normal NLP in \mathbf{x}
- Alternative notions to handle uncertainty:
 - Parametric Optimization
 - **Stochastic Programming**
 - **Robust Optimization**

Alternatives to Parametric Optimization

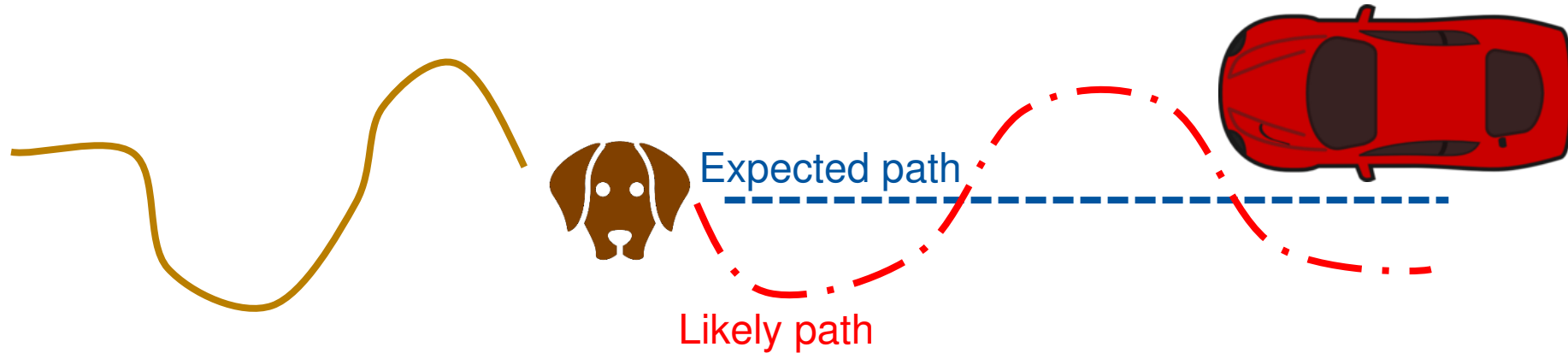
Different notions to account for uncertainty:

- **Stochastic approaches** consider probability measures over the set of possible uncertainty realizations
- **Robust approaches** consider worst-case uncertainty realization

	Stochastic	Robust
Assumption	Known probability for \mathbf{y}	\mathbf{y} bounded, $\mathbf{y} \in Y$
Objective	e.g., $\min_{\mathbf{x}} \mathbb{E}_Y(f(\mathbf{x}, \mathbf{y}))$	$\min_{\mathbf{x}} \max_{\mathbf{y} \in Y} f(\mathbf{x}, \mathbf{y})$
	Optimality in probabilistic sense	Optimality for the worst case
Constraint	e.g., $\mathbb{P}_Y(\mathbf{c}(\mathbf{x}, \mathbf{y}) \leq \mathbf{0}) \geq \alpha, \alpha \in (0,1)$	$\mathbf{c}(\mathbf{x}, \mathbf{y}) \leq \mathbf{0}, \forall \mathbf{y} \in Y$
	Chance for feasibility	Guaranteed feasibility

- Compared to neglecting uncertainties, both approaches result in more reliable and conservative solutions

Example: Optimization under Uncertainty (1)



Inspired by Sam L. Savage, Stanford University

Nominal optimization:

- Consider expected path, **drive straight**. Likely result: 

Stochastic approach:

- Consider likely paths, **drive to the side and slow down**. Likely result: 

Robust approach:

- Consider all possible paths, **stop the car**. Guaranteed result: 

Example: Optimization under Uncertainty (2)



Minimize time at airport ($T@A$), missed flight \Rightarrow 3h wait

- Nominal optimization:
 - Leave 58.5 min before flight: avg. $T@A$: 90min, 50% missed flights
- Stochastic approach:
 - Leave 60 min before flight: avg. $T@A$: 19.5min, 10% missed flights
- Robust approach:
 - Leave 120 min before flight: avg. $T@A$: 61.5min, 0% missed flights

Check Yourself

- What problems can arise if uncertainty is neglected in the solution of optimization problems?
- Name the basic notions of how uncertainty can be addressed. How do the approaches relate to each other?



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Two-stage stochastic optimization

Two-Stage Stochastic Programming

Consider the following problem description:

- $\mathbf{y} \in Y$ is an **uncertain parameter** with some probability distribution.
- $\mathbf{x} \in X$ is a **decision variable** that must be fixed **before** the uncertainty is realized with constraints $\mathbf{c}^U(\mathbf{x}) \leq \mathbf{0}$.
- $\mathbf{z} \in Z$ is a **decision variable** that can be fixed **after** the uncertainty is realized with constraints $\mathbf{c}^L(\mathbf{x}, \mathbf{y}, \mathbf{z}) \leq \mathbf{0}$.
- An objective function $f : X \times Y \times Z \rightarrow \mathbb{R}$ is to be minimized.
- f is separable such that $f(\mathbf{x}, \mathbf{y}, \mathbf{z}) = f^U(\mathbf{x}) + f^L(\mathbf{y}, \mathbf{z})$.
- The corresponding two-stage stochastic program can be written as

$$\begin{aligned} \min_{\mathbf{x}} \quad & f^U(\mathbf{x}) + \mathbb{E}_Y(F(\mathbf{x}, \mathbf{y})) & F(\mathbf{x}, \mathbf{y}) = \min_{\mathbf{z}} \quad & f^L(\mathbf{y}, \mathbf{z}) & \text{(ST2)} \\ \text{s. t.} \quad & \mathbf{c}^U(\mathbf{x}) \leq \mathbf{0} & & \text{s. t.} \quad & \mathbf{c}^L(\mathbf{x}, \mathbf{y}, \mathbf{z}) \leq \mathbf{0} & \text{(ST1)} \end{aligned}$$

- (ST1) determines the **first-stage** decision process before the uncertainty is realized.
- (ST2) describes the **second-stage** decision process for a specific uncertainty scenario.

Two-Stage Stochastic Programming Example: Plant design

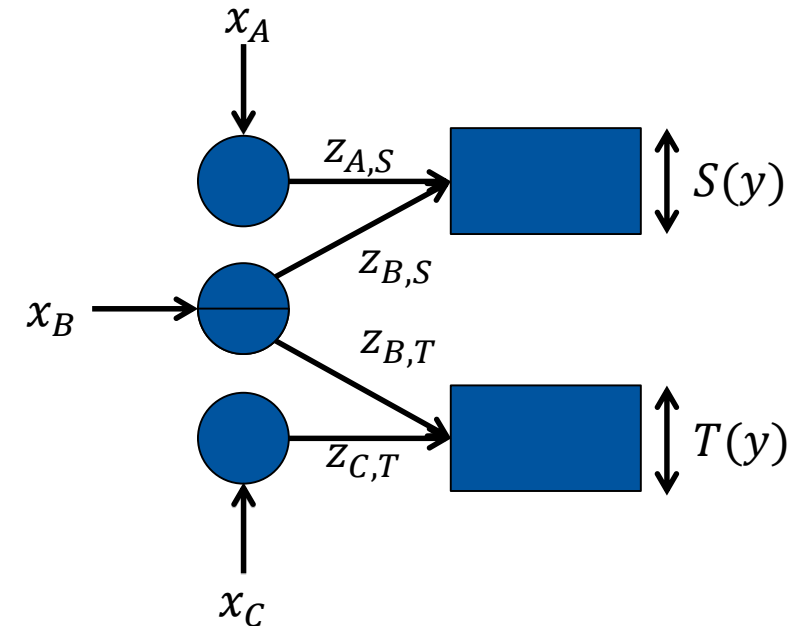
Determine **sizing** (x_A, x_B, x_C) of different production tracks to **maximize profit** under **uncertain demand** $(S(y), T(y))$.

1. Decide on investment into tracks A, B, C : (x_A, x_B, x_C) .
2. Realize uncertainty y and obtain demand: $S(y)$ and $T(y)$.
3. Decide on production of products S and T from tracks A, B, C : $(z_{A,S}, z_{B,S}, z_{B,T}, z_{C,T})$.

$$\min_x d_A x_A + d_B x_B + d_C x_C + \mathbb{E}_Y(F(x, y))$$

$$\begin{aligned} \text{s. t. } & x_A + x_B + x_C \leq x_{\max} \\ & x_A, x_B, x_C \geq 0 \end{aligned}$$

$$\begin{aligned} F(x, y) = \min_z & -d_S(z_{A,S} + z_{B,S}) - d_T(z_{B,T} + z_{C,T}) \\ \text{s. t. } & z_{A,S} \leq x_A & z_{A,S} + z_{B,S} \leq S(y) \\ & z_{B,S} + z_{B,T} \leq x_B & z_{B,T} + z_{C,T} \leq T(y) \\ & z_{C,T} \leq x_C \end{aligned}$$



A Single-Stage Reformulation

$$\begin{aligned} \min_{\mathbf{x}} \quad & f^U(\mathbf{x}) + \mathbb{E}_Y(F(\mathbf{x}, \mathbf{y})) \\ \text{s. t.} \quad & \mathbf{c}^U(\mathbf{x}) \leq \mathbf{0} \end{aligned} \quad (\text{ST1})$$

$$\begin{aligned} F(\mathbf{x}, \mathbf{y}) = \min_{\mathbf{z}} \quad & f^L(\mathbf{y}, \mathbf{z}) \\ \text{s. t.} \quad & \mathbf{c}^L(\mathbf{x}, \mathbf{y}, \mathbf{z}) \leq \mathbf{0} \end{aligned} \quad (\text{ST2})$$

- If \mathbf{y} has a continuous or large discrete distribution, $\mathbb{E}_Y(F(\mathbf{x}, \mathbf{y}))$ is difficult if not impossible to evaluate exactly.
- The distribution can be approximated by finitely many scenarios $s \in \mathcal{S}$ with $\mathbf{y} = \mathbf{y}_s > \mathbf{0}$ and probability of occurrence P_s such that $\sum_{s \in \mathcal{S}} P_s = 1$

- Then, the two-stage stochastic program can be approximated by the single-stage program

$$\begin{aligned} \min_{\mathbf{x}, \mathbf{z}_s} \quad & f^U(\mathbf{x}) + \sum_{s \in \mathcal{S}} P_s \cdot f^L(\mathbf{y}_s, \mathbf{z}_s) \\ \text{s. t.} \quad & \mathbf{c}^U(\mathbf{x}) \leq \mathbf{0} \\ & \mathbf{c}^L(\mathbf{x}, \mathbf{y}_s, \mathbf{z}_s) \leq \mathbf{0}, \forall s \in \mathcal{S} \end{aligned}$$

- The resulting problem is potentially very large and challenging to solve.
- More efficient approaches have been proposed in literature.

Check Yourself

- What is the essential feature of a two-stage stochastic program?
- When can a two-stage stochastic program be reformulated exactly into a single-stage program?



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Introduction to Semi Infinite Programs

Semi-Infinite Optimization Problems

What does the following mean?

$$\begin{aligned} \min_{x \in X} f(x) \\ \text{s.t. } c(x, y) \leq 0, \forall y \in Y \end{aligned} \quad (\text{SIP})$$

- A point $x \in X$ is feasible if and only if $c(x, y) \leq 0$ holds for all possible values of $y \in Y$.
- If $|Y| = \infty$, we speak of a **semi-infinite** optimization problem (SIP)
- For simplicity of notation we take a single SIP constraint, and a single uncertain variable $y \in Y \subset R$
- SIP has finitely many variables and **infinitely many constraints**
- SIPs date back to at least 1960s
- SIPs are useful for worst-case optimization="robust optimization"
- Many generalizations exist, including to $Y(x)$, existence constraints
- It is important special case of hierarchical problems (bilevel, trilevel)

Semi-Infinite Optimization Problems: A Useful Reformulation

$$\begin{aligned} \min_{x \in X} f(x) \\ \text{s.t. } c(x, y) \leq 0, \forall y \in Y \end{aligned} \quad (\text{SIP})$$

- If $\max_{y \in Y} c(x, y)$ exists, we can rewrite (SIP) as

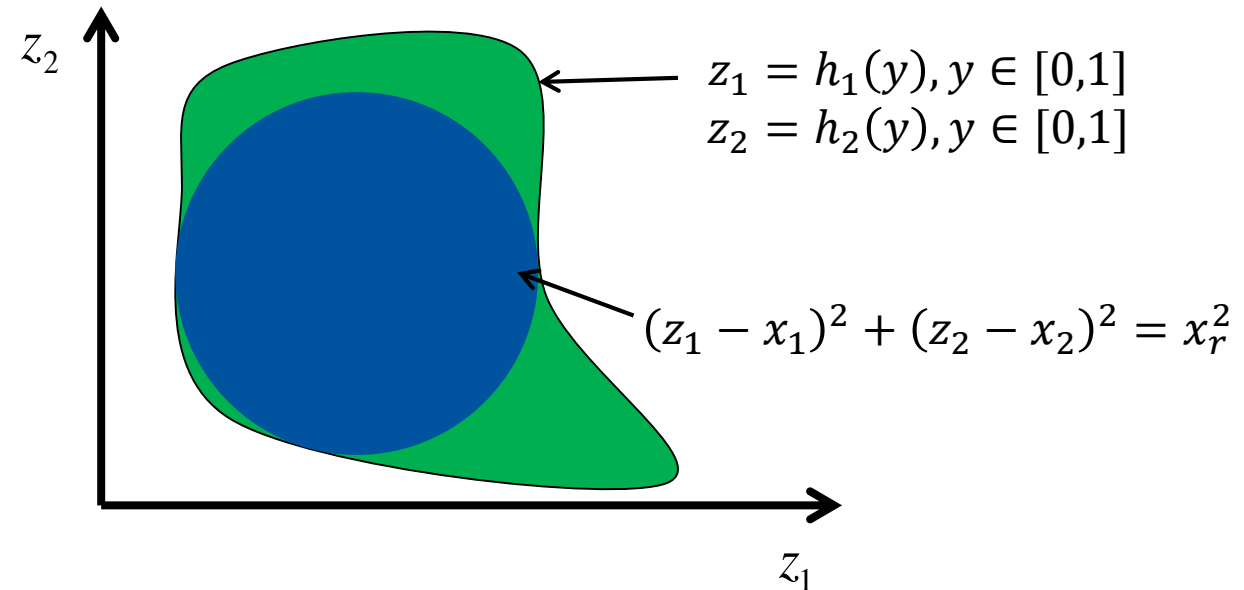
$$\begin{aligned} \min_{x \in X} f(x) \\ \text{s.t. } 0 \geq \max_{y \in Y} c(x, y) \end{aligned}$$

- We call $\max_{y \in Y} c(x, y)$ the **lower-level problem** (LLP)
- A point $x \in X$ is feasible if and only if the **global solution** to LLP is ≤ 0

Design Centering as Semi-Infinite Program

- Fit the largest **Shape A** that fits into **Shape B**
- Example: circle
 - Variables: center coordinates & radius $\mathbf{x} = (x_1, x_2, x_r)$
 - Objective: maximal radius $\Rightarrow \max_{\mathbf{x}} x_r$
 - Constraint: **Shape A** shall fit into **Shape B**

$$\Rightarrow (h_1(y) - x_1)^2 + (h_2(y) - x_2)^2 \geq x_r^2, \forall y \in [0,1]$$



- Real-life applications:
 - Diamond cutting: Stein, Optimization with Multivalued Mappings: Theory, Applications and Algorithms, 2006
 - Model reduction: Oluwole et. al, Combustion and Flame, 2006

Example: Path Constraints in Dynamic Optimization

- Dynamic optimization: infinite # variables
 - Handled by control-vector parametrization
- States uniquely determined as function of time by model
- Path constraints hold for $\forall t \rightarrow$ semi-infinite problem
- Standard algorithms & solvers enforce the path constraints for a finite # times:
 - $[0, t_f]$ is discretized to $\{t_1, t_2, \dots, t_N\} \subset [0, t_f]$
 - Relaxation that may result in violations
- Motivates the use of SIP techniques^{1,2}

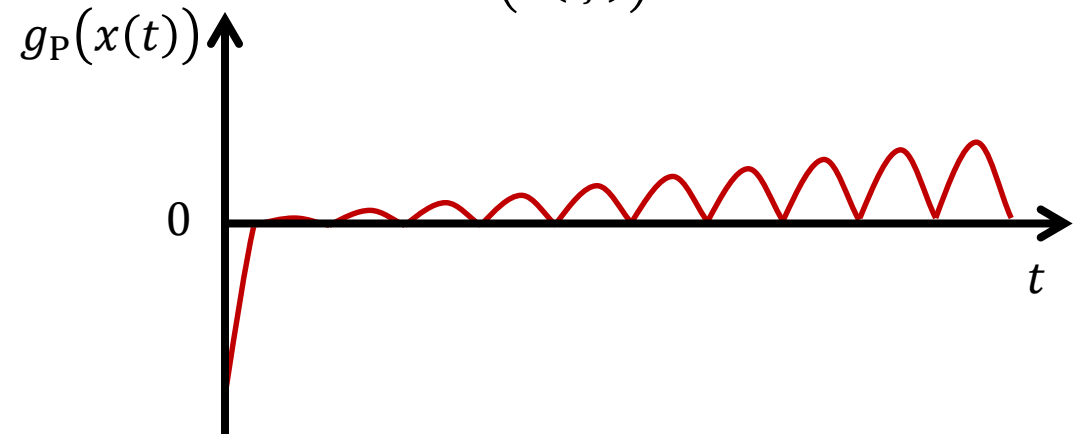
$$\min_{x(\cdot), u(\cdot)} \Phi(x(t_f))$$

$$\text{s.t.} \quad \dot{x}(t) = f(x(t), u(t)), t \in [t_0, t_f]$$

$$x(t_0) = x_0$$

$$g_P(x(t), u(t)) \leq 0 \quad \forall t \in [t_0, t_f]$$

$$g_T(x(t_f)) \leq 0$$



Example: Optimal Operation Under Parametric Uncertainty

$$\min_{x(\cdot), u(\cdot)} \Phi(x(t_f))$$

$$\text{s.t.} \quad \dot{x}(t) = f(x(t), u(t), y), t \in [t_0, t_f]$$

$$x(t_0) = x_0(y)$$

$$g_P(x(t), u(t), y) \leq \mathbf{0} \quad \forall t \in [t_0, t_f]$$

$$\min_{u(\cdot)} \Phi(x(t_f))$$

$$\text{s.t.} \quad g_P(x(t), u(t), y) \leq \mathbf{0} \quad \forall t \in [t_0, t_f]$$

Where $x(\cdot)$ is the solution of

$$\dot{x}(t) = f(x(t), u(t), y), t \in [t_0, t_f]$$

$$x(t_0) = x_0(y)$$

- Recall sequential interpretation
- Worst-case formulation by introducing parametric solution of states

$$\min_{u(\cdot)} \Phi(x(t_f))$$

$$g_P(x_y(t), u(t), y) \leq \mathbf{0} \quad \forall y, \forall t \in [t_0, t_f]$$

Where $x(\cdot)$ is the solution of nominal $y = \bar{y}$

$$\dot{x}(t) = f(x(t), u(t), \bar{y}), t \in [t_0, t_f]$$

$$x(t_0) = x_0(\bar{y})$$

and $x_y(\cdot)$ is the solution for given y

$$\dot{x}(t) = f(x(t), u(t), y), t \in [t_0, t_f]$$

$$x(t_0) = x_0(y)$$

Check Yourself

- For which kind of constraint functions is an SIP particularly difficult to solve? Why is this the case?
- What are applications of SIP?



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Basic solution methods for Semi Infinite Programs

Semi-Infinite Optimization Problems: A Useful Reformulation

$$\begin{aligned} \min_{x \in X} f(x) \\ \text{s.t. } c(x, y) \leq 0, \forall y \in Y \end{aligned} \quad (\text{SIP})$$

- If $\max_{y \in Y} c(x, y)$ exists, we can rewrite (SIP) as

$$\begin{aligned} \min_{x \in X} f(x) \\ \text{s.t. } 0 \geq \max_{y \in Y} c(x, y) \end{aligned}$$

- We call $\max_{y \in Y} c(x, y)$ the **lower-level problem** (LLP)
- A point $x \in X$ is feasible if and only if the **global solution** to LLP is ≤ 0

Intuitive Solution Approach as Nested Problem?

- Treat semi-infinite constraint as black box function

$$c^*(\mathbf{x}) = \max_{y \in Y} c(\mathbf{x}, y)$$

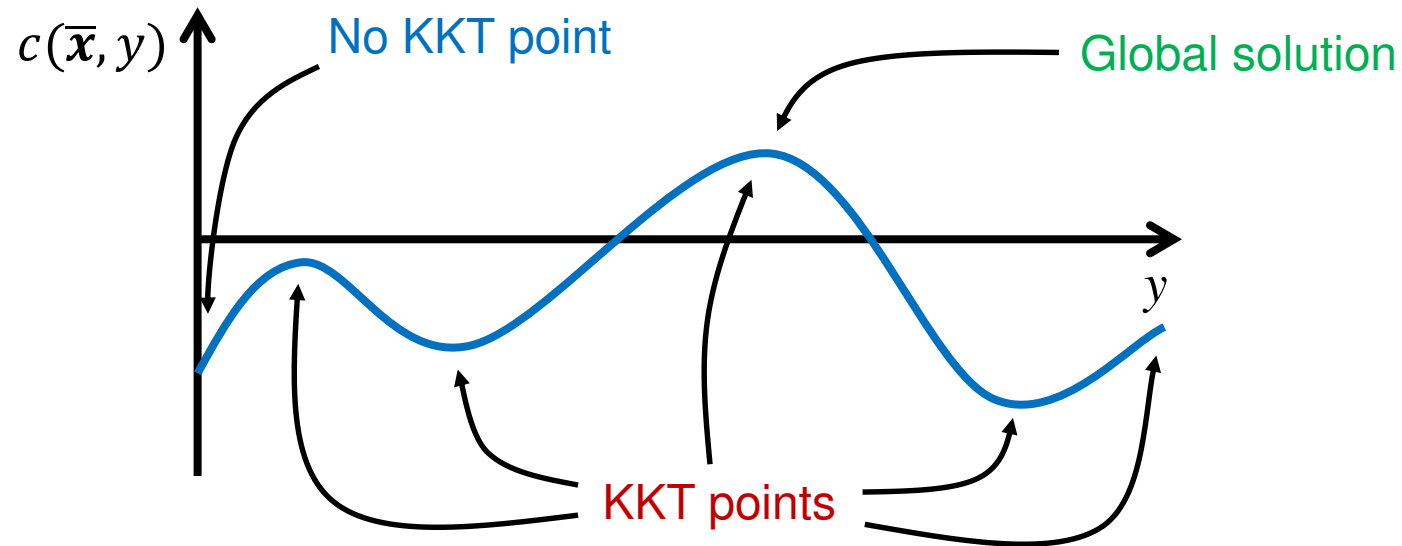
- Use local solver to solve

$$\begin{array}{ll} \min_{\mathbf{x} \in X} & f(\mathbf{x}) \\ \text{s.t.} & c^*(\mathbf{x}) \leq 0 \end{array}$$

- For any particular $\mathbf{x} \in X$ obtain $c^*(\mathbf{x})$ by solving (LLP) globally
- Problems
 - Gradient-based methods are not applicable: $c^*(\cdot)$ is not differentiable
 - Gradient-free local solvers evaluate many points $\mathbf{x} \in X$ requiring many global solutions of (LLP)
 - Parametric solution of $c^*(\cdot)$ is extremely expensive
- Nonsmooth local solvers of interest

Intuitive Solution Approach using Optimality Conditions of LLP?

- Replace the lower-level problem with its KKT conditions
- The resulting problem is finite and can be solved by NLP methods
 - Recall that the KKT conditions have nonsmoothness
- In general **wrong** since KKT conditions are **not sufficient for the global** maximization of LLP



- Relaxation of the SIP: resulting point is not necessarily feasible in the SIP

Solution Methods of SIP: Local Reduction and Discretization

1. Local reduction

- For a point $\bar{x} \in X$, find all KKT points of the LLP
- Track all KKT points
- The resulting problem is finite and can be solved by NLP methods

2. Discretization

- Replace Y with a finite discretization $Y^D \subset Y$
- The resulting problem is a finite approximation of the SIP and can be solved by NLP methods
- The resulting problem is a relaxation
- The discretization Y^D is populated to better approximate the SIP
- The relaxation can be tightened by using also KKT conditions of LLP

Check Yourself

- For which kind of constraint functions is an SIP particularly difficult to solve? Why is this the case?
- Why not solve the nested problem directly?
- Name the basic approaches to the solution of SIPs.
- How does the discretized version of an SIP relate to the original problem?



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Deterministic global solution of Semi Infinite Programs

Semi-Infinite Optimization Problems: A Useful Reformulation

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- If $\max_{y \in Y} c(x, y)$ exists, we can rewrite (SIP) as

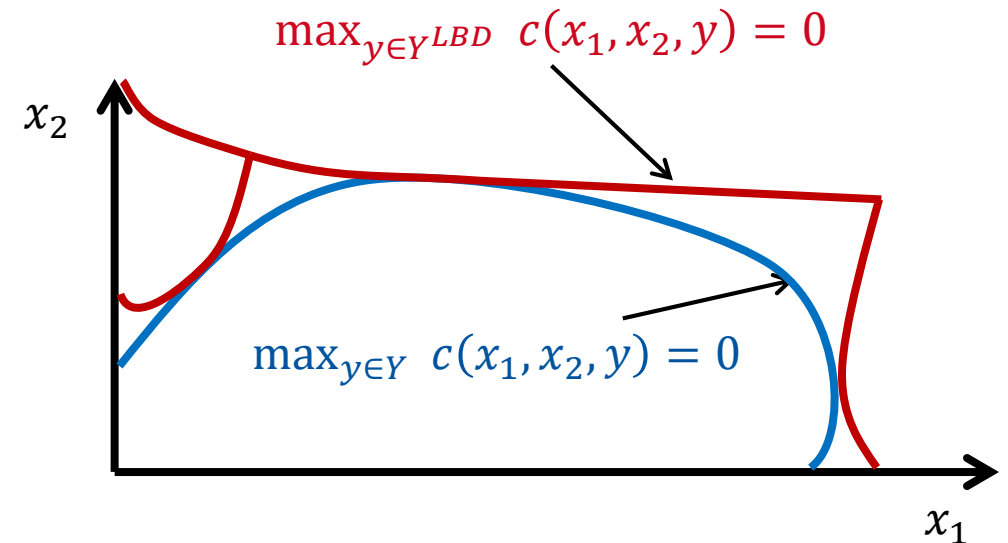
$$\begin{aligned} \min_{x \in X} \quad & f(x) \\ \text{s.t.} \quad & 0 \geq \max_{y \in Y} c(x, y) \end{aligned}$$

- We call $\max_{y \in Y} c(x, y)$ the **lower-level problem** (LLP)
- Goal: solve the SIP globally using discretization-based lower & upper bounding procedures
 - up to optimality gap $f^{UBD} - f^{LBD} \leq \varepsilon^f$
 - Assuming an SIP Slater point, continuous functions and compact X, Y

Discretization Approach from Blankenship & Falk

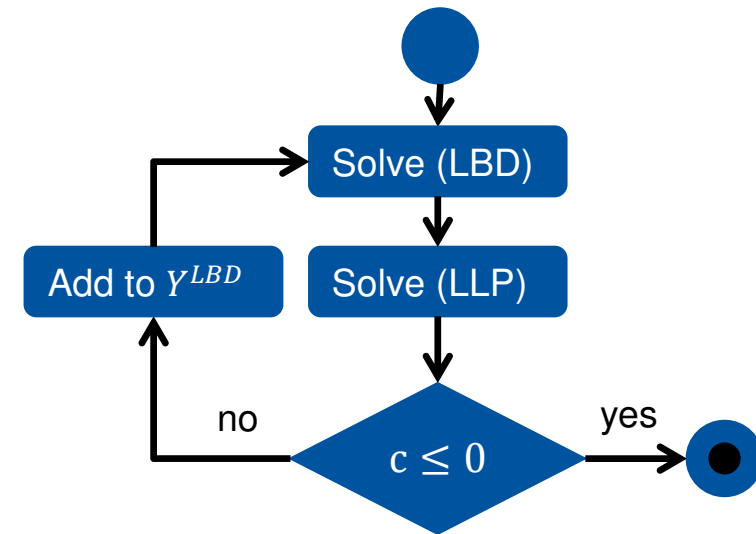
- Approximate Y with $Y^{LBD} \subset Y$, $|Y^{LBD}| < \infty$ to lower bounding problem (LBD)
 - The LLP is restricted and thus the SIP relaxed
 - Outer approximation
- Solve (LBD) and obtain \bar{x}
- Solve (LLP) **globally** for $x = \bar{x}$ to alternatively obtain
 - $c^*(\bar{x}) \leq 0$, \bar{x} is SIP feasible
 - \bar{y} : $c(\bar{x}, \bar{y}) > 0$

$$\begin{aligned} \min_{x \in X} \quad & f(x) \\ \text{s.t.} \quad & c(x, y) \leq 0, \forall y \in Y^{LBD} \end{aligned} \quad (\text{LBD})$$



Lower Bounding Procedure

- (LBD) is a relaxation of (SIP) for any $Y^{LBD} \subset Y$
- The global solution of (LBD) gives a lower bound on the globally optimal objective value of (SIP)
- Algorithm:
 1. Initialize $Y^{LBD} \subset Y, f^{LBD} \leftarrow -\infty$
 2. Solve (LBD) globally to obtain \bar{x}
 - Set $f^{LBD} \leftarrow f(\bar{x})$
 3. Solve (LLP) globally for \bar{x} to obtain \bar{y}
 - If $c(\bar{x}, \bar{y}) \leq 0$ then set $x^* \leftarrow \bar{x}$ and terminate (x^* is feasible)
 - Else set $Y^{LBD} \leftarrow \{\bar{y}\} \cup Y^{LBD}$ and go to step 2
- The algorithm converges to the global solution to (SIP) in the limit
- Finite generation of a feasible point is not guaranteed



Discretization with Guaranteed Feasibility

- Recall (LBD)

$$\begin{aligned} \min_{x \in X} \quad & f(x) \\ \text{s.t.} \quad & c(x, y) \leq 0, \forall y \in Y^{LBD} \end{aligned} \quad (\text{LBD})$$

- Restrict ($\varepsilon > 0$) the right-hand-side of the constraint to obtain the upper bounding problem (UBD)¹

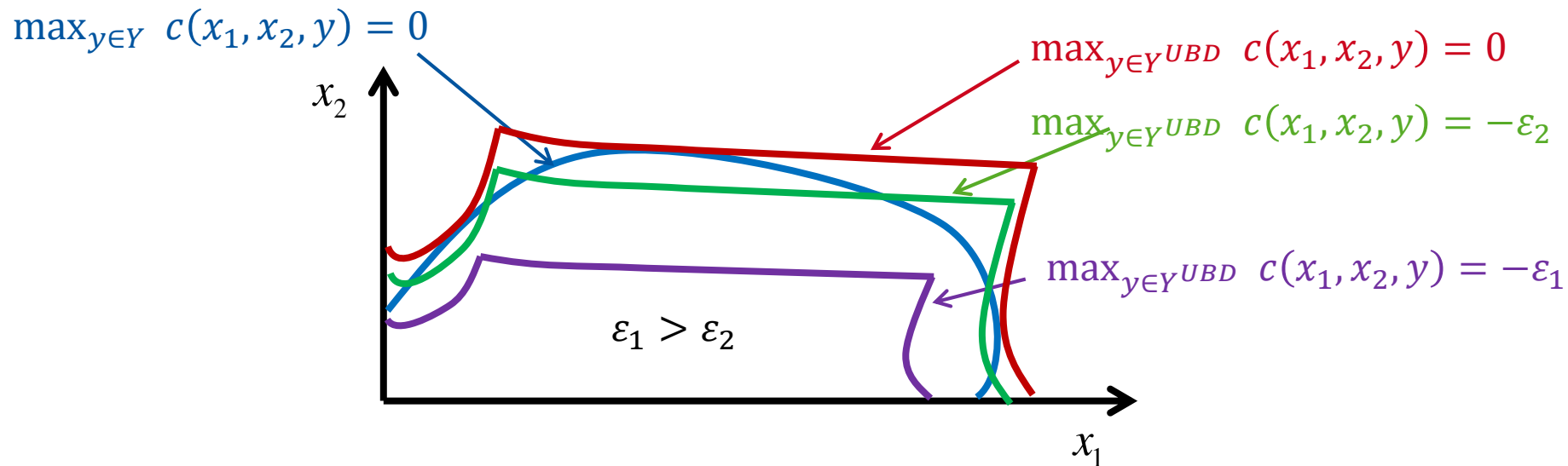
$$\begin{aligned} \min_{x \in X} \quad & f(x) \\ \text{s.t.} \quad & c(x, y) \leq -\varepsilon, \forall y \in Y^{UBD} \end{aligned} \quad (\text{UBD})$$

- (UBD) is neither a relaxation nor a restriction of (SIP)

Upper Bounding Problem

- (UBD) is neither a relaxation nor a restriction of (SIP)

$$\begin{aligned} \min_{x \in X} \quad & f(x) \\ \text{s.t.} \quad & c(x, y) \leq -\varepsilon, \forall y \in Y^{UBD} \end{aligned} \quad (\text{UBD})$$



- By populating Y^{UBD} , the relaxation is made tighter \rightarrow better approximate (SIP)
- By reducing $\varepsilon > 0$, the restriction is less severe \rightarrow better approximate (SIP)

Simplified Upper Bounding Procedure

- Algorithm:

1. Initialize $Y^{UBD} \subset Y, f^{UBD} \leftarrow \infty, \varepsilon > 0, r > 1$

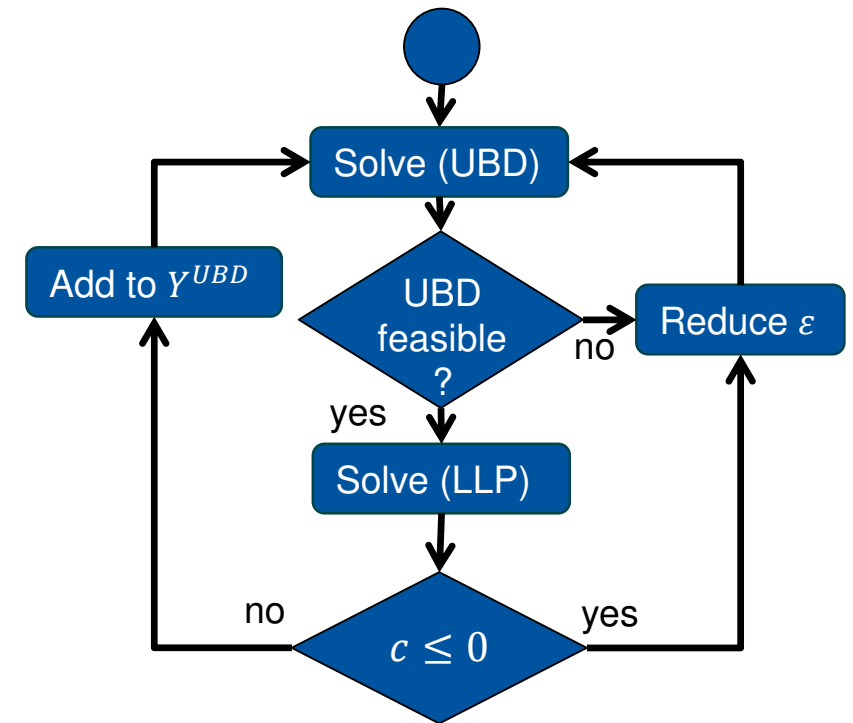
2. Solve (UBD) globally to obtain \bar{x}

3. Solve (LLP) globally for \bar{x} to obtain \bar{y}

- If $c(\bar{x}, \bar{y}) \leq 0$ then set $\mathbf{x}^* \leftarrow \bar{x}$ and $\varepsilon \leftarrow \varepsilon/r$
- Else set $Y^D \leftarrow \{\bar{y}\} \cup Y^{UBD}$
- Go to step 2

- Finite convergence to global optimum if SIP-Slater point exists

- If for $\varepsilon^f > 0, \exists \mathbf{x}^S \in X: c(\mathbf{x}^S, y) < 0, \forall y \in Y, f(\mathbf{x}^S) \leq f^* + \varepsilon^f$ the procedure finitely produces an ε^f -optimal, feasible point



Check Yourself

- How does the discretized version of an SIP relate to the original problem?
- Describe a global solution method for SIP



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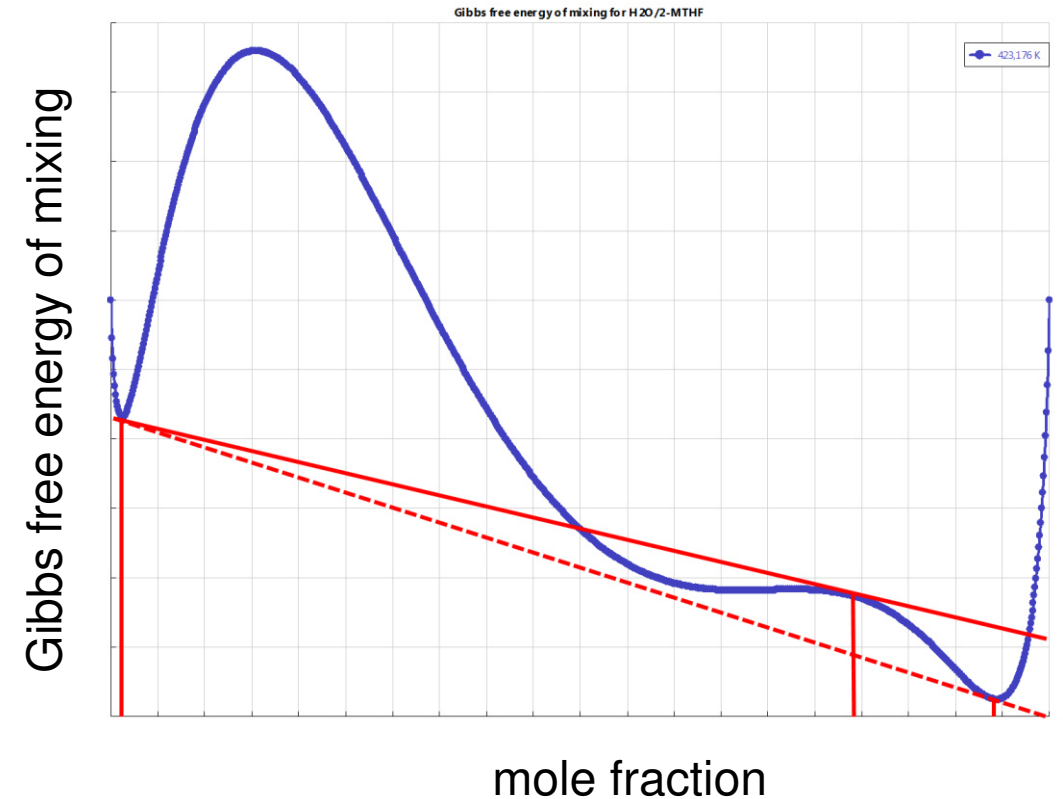
Parameter estimation in thermodynamics as bilevel/SIP

Parameter Estimation for Thermodynamic Property Models

- Accurate prediction of phase separation behavior is needed, e.g., in separation design.
- Excess-Gibbs free energy models are used for liquid phases $G(T, P, \mathbf{x}, \mathbf{q})$
 - NRTL, UNIQUAC, Wilson, ... :
- EOS models are used for vapor phases
 - No adjustment of parameters for binary interaction
- Binary interaction parameters \mathbf{q} estimated via equilibrium experiments
- Parameters must be fit for model equations in equilibrium state.
 - Necessary equilibrium criterion: Isopotential (equality of chemical potential between phases)
 - Necessary and sufficient equilibrium criterion: $\min_{\mathbf{x}} G(T, P, \mathbf{x}, \mathbf{q})$
 - Specialized criteria exist: Gibbs tangent plane (Michelsen 1980) supporting hyperplane (Mitsos & Barton AIChEJ 2007)

Parameter Estimation: Classical approach

- Measure phase compositions $x^{m,i,k}$ at different temperatures T^i and constant pressure P .
- Minimize prediction-measurement discrepancy
- Enforce isopotential for thermodynamic equilibrium
- $\min_q LS(x^{p,i,k} - x^{m,i,k})$
 $\mu_1^1(T^i, P, x^{p,i,1}, \mathbf{q}) = \mu_1^k(T^i, P, x^{p,i,k}, \mathbf{q}) \quad \forall i, k$
 $\mu_2^1(T^i, P, x^{p,i,1}, \mathbf{q}) = \mu_2^k(T^i, P, x^{p,i,k}, \mathbf{q}) \quad \forall i, k$
- Problem: isopotential is only **necessary**, not **sufficient**.



Glass, et al., Mitsos, *Fluid Phase Equilibria* 433 (2017): 212-225.

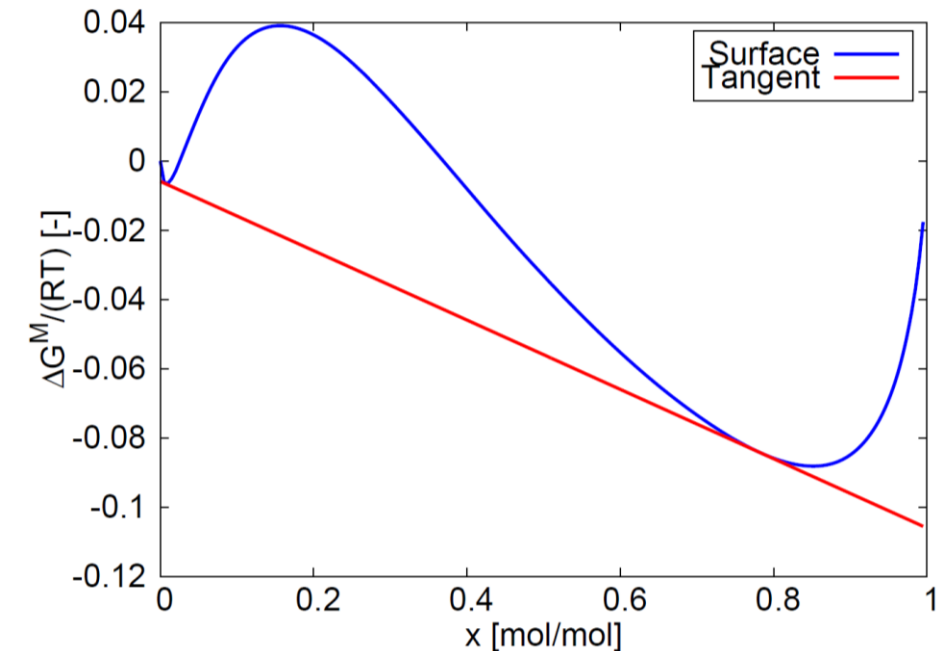
Parameter Estimation: Bilevel Formulation

- Developed by Mitsos, Bolas & Barton CES 2009
 - Extended by Glass & Mitsos
 - Implemented in BOARPET

- Introduce tangent plane criterion

$$\Delta g(T^i, P^i, x^{p,i,1}, \mathbf{q}) + \left. \frac{\delta \Delta g}{\delta x} \right|_{T^i, P^i, x^{p,i,1}, \mathbf{q}} (x - x^{p,i,1}) \leq \Delta g(T^i, P^i, x, \mathbf{q}) \quad \forall x \in [0,1]$$

- Additional constraints:
 - correct number of phase splits
 - correct number of phases for each split



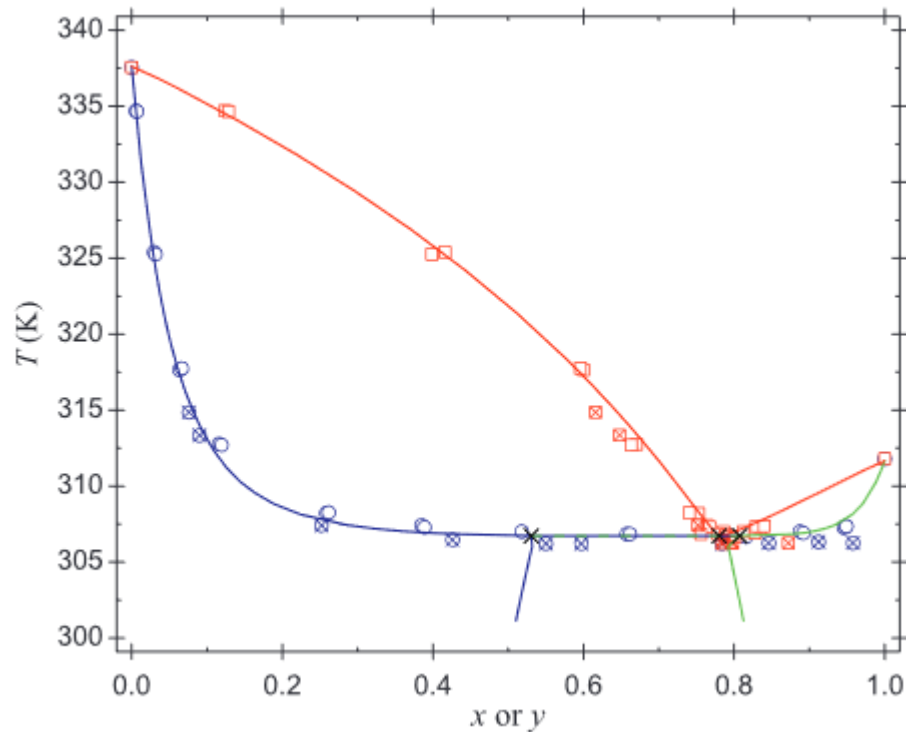
Glass, et int., Mitsos, *Fluid Phase Equilibria* 433 (2017): 212-225.

<https://www.avt.rwth-aachen.de/cms/AVT/Forschung/Software/~kvkz/BOARPET/>

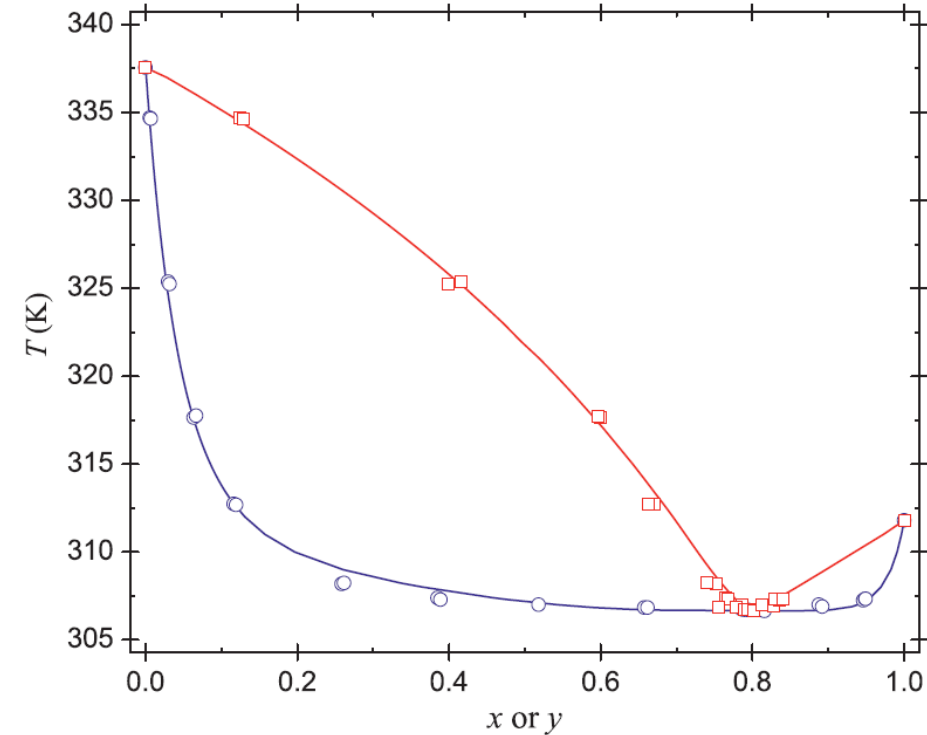
Parameter Estimation: VLE of 2-methyl-2-butene – methanol

- Experimental data shows homogeneous azeotrope

$\gamma - \phi$ method predicts heterogeneous azeotrope (**spurious LLE split**).



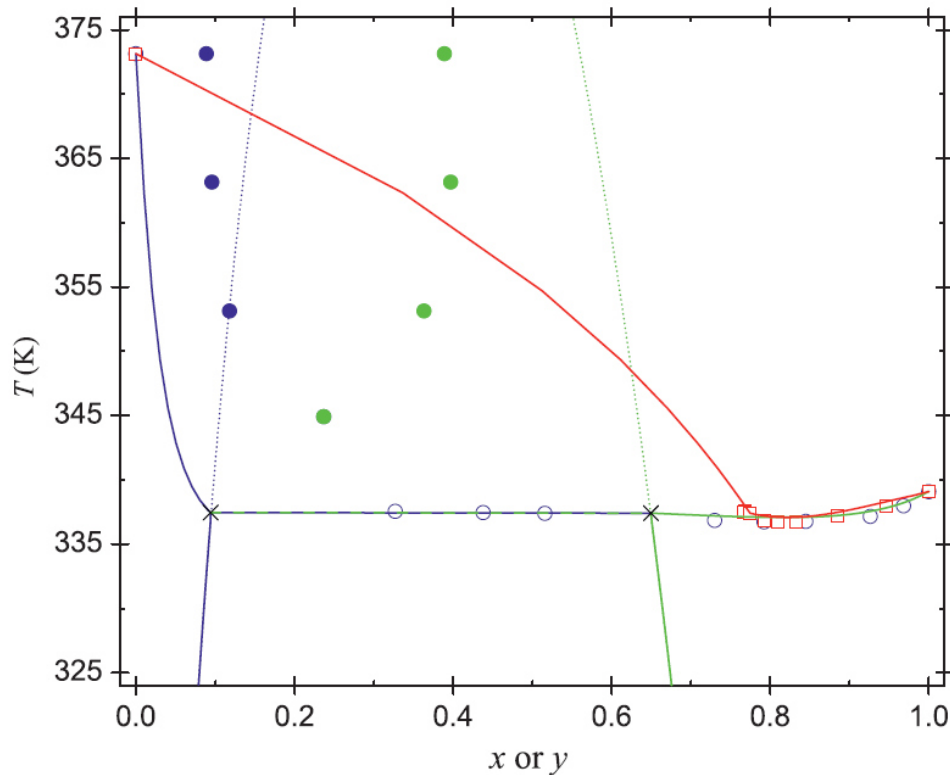
Bilevel formulation gives qualitatively and quantitatively correct fit.



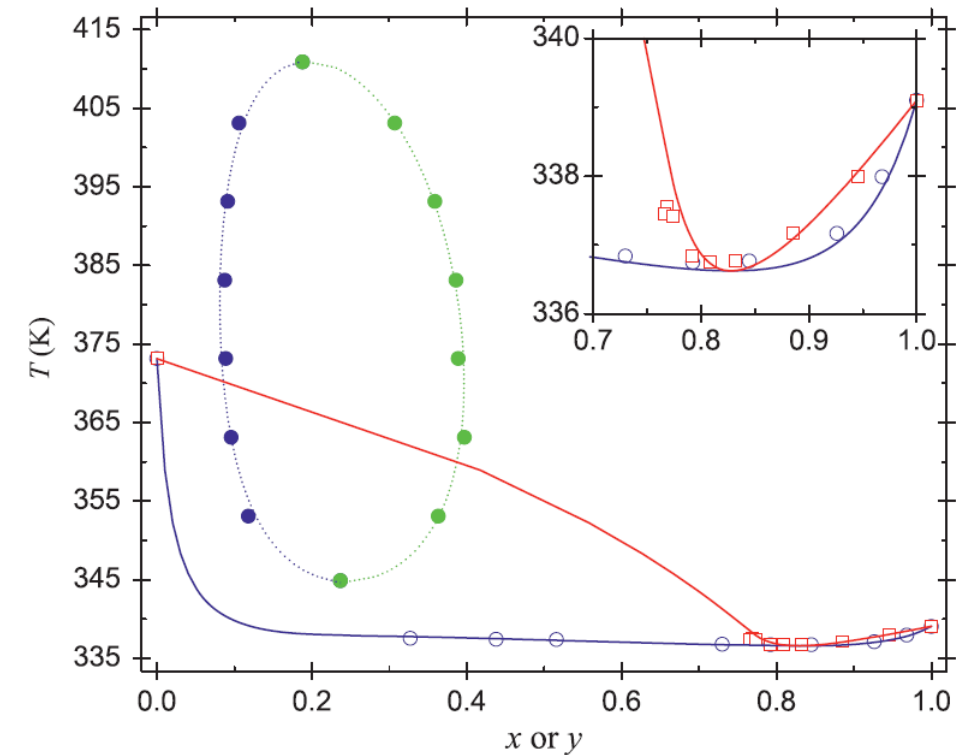
Parameter Estimation: VLE of tetrahydrofuran-water

- Experimental data shows homogeneous azeotrope, LCST, and UCST

$\gamma - \phi$ method predicts **spurious liquid phase split**.
LCST is missed



Proposed formulation gives qualitatively and quantitatively correct fit.



Check Yourself

- How does parameter estimation in thermodynamics work?