

# **Applied Numerical Optimization**

Prof. Alexander Mitsos, Ph.D.

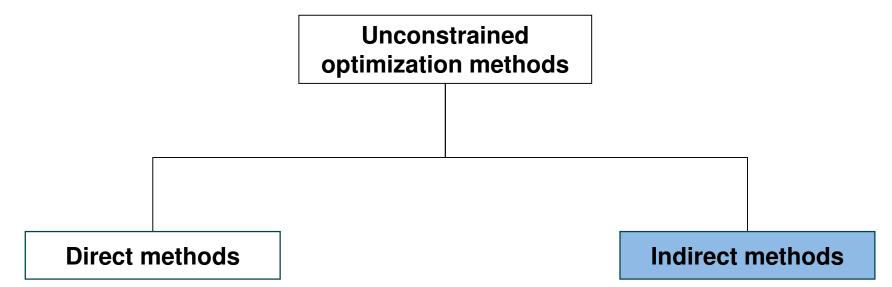
Basic solution methods for unconstrained problems





# **Solution Methods for Unconstrained Optimization**

 $\min_{\mathbf{x}\in R^n}f(\mathbf{x})$ 





# **Indirect Methods – Concept**

First-order necessary conditions

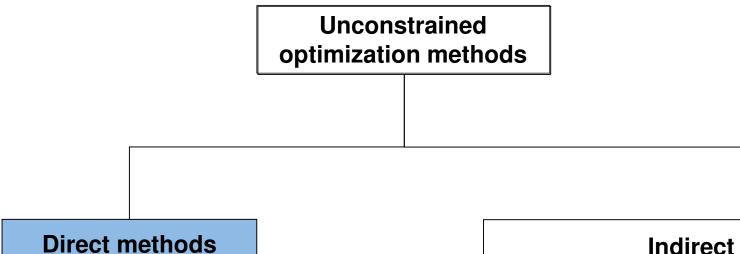
$$\nabla f(x) = \mathbf{0} \Leftrightarrow \begin{cases} \left. \frac{\partial f}{\partial x_1} \right|_x = 0 = g_1(x) & \text{nonlinear system of equations} \\ \left. \frac{\partial f}{\partial x_2} \right|_x = 0 = g_2(x) \Leftrightarrow \mathbf{g}(x) = \mathbf{0} \\ \vdots \\ \left. \frac{\partial f}{\partial x_1} \right|_x = 0 = g_n(x) \end{cases}$$

- The optimal solution is found by solving the system of equations analytically or numerically (e.g., by Newton's method).
- Differentiation and solution of the system of equations is challenging for complex problems!





# **Solution Methods for Unconstrained Optimization**



#### **Indirect methods**

Optimal solution is found by solving the system of equations (optimality conditions):

$$\nabla f(x) = 0$$

analytically or numerically

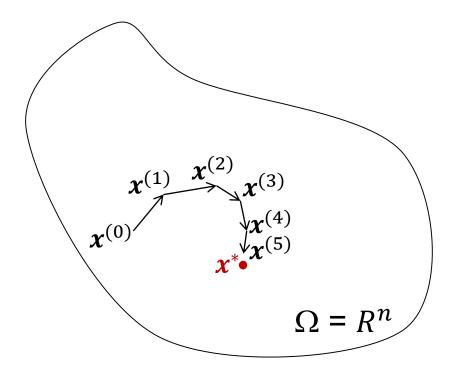




# **Direct Methods – Concept**

Idea: Construct a convergent sequence of  $\{x^{(k)}\}_{k=1}^{\infty}$ , which fulfills the following conditions:

$$\exists \overline{k} \ge 0$$
:  $f(x^{(k+1)}) < f(x^{(k)}) \forall k > \overline{k}$  and  $\lim_{k \to \infty} x^{(k)} = x^* \in R^n$ 





# **Definition: Rate of Convergence**

Idea: Construct a convergent sequence of  $\{x^{(k)}\}_{k=1}^{\infty}$ , which fulfills the following conditions:

$$\exists \bar{k} \ge 0$$
:  $f(x^{(k+1)}) < f(x^{(k)}) \forall k > \bar{k}$  and  $\lim_{k \to \infty} x^{(k)} = x^* \in R^n$ 

#### Rate of convergence:

• Linear: if there exists a constant  $C \in (0,1)$ , such that for sufficiently large k:

$$||x^{(k+1)} - x^*|| \le C||x^{(k)} - x^*||$$

• Order p (often p = 2): if there exists a constant M > 0, such that

$$||x^{(k+1)} - x^*|| \le M ||x^{(k)} - x^*||^p$$

• Superlinear: if there exists a sequence  $c_k$  converging to zero, i.e.,  $\lim_{k\to\infty}c_k=0$ , such that

$$\|x^{(k+1)} - x^*\| \le c_k \|x^{(k)} - x^*\|$$





# **Solution Methods for Unconstrained Optimization**

Unconstrained optimization methods

#### **Direct methods**

Optimal solution is found by directly improving the objective function via iterative descent.

#### **Indirect methods**

Optimal solution is found by solving the system of equations (optimality conditions):

$$\nabla f(x) = 0$$

analytically or numerically.





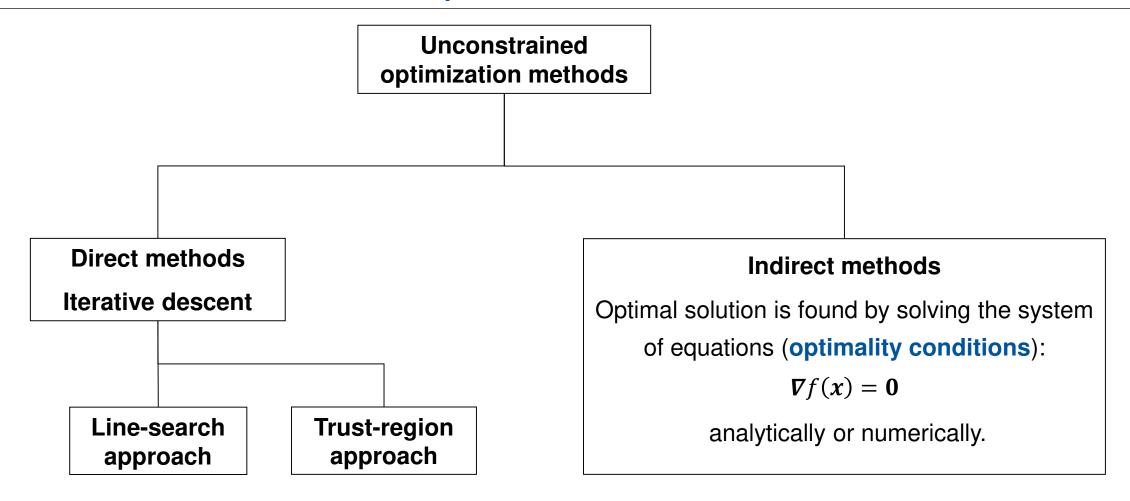
#### **Direct vs Indirect: Nomenclature not consistent in Literature**

- Throughout class we use "direct" and "indirect":
  - "indirect methods": 1. set up optimality conditions and then 2. try to solve the system of equations (or equations and inequalities)
  - "direct methods": directly aim to improve objective function (or objective function and constraints). These methods hope to converge to optimality conditions.
- In the literature there are many alternative uses of the word, including
  - exactly the opposite than ours
  - "direct": without the use of derivatives, "indirect": using derivatives
  - only in the context of dynamic optimization problems:
    - "direct": first convert to nonlinear program
    - "indirect": first set up optimality conditions
  - only in the context of constrained problems
    - "direct": only feasible iterates
    - "indirect": infeasible iterates are allowed





# **Solution Methods for Unconstrained Optimization**





#### **Check Yourself**

- What are direct vs indirect methods?
- Which direct methods did we learn?
- Which convergence rates exist? Why is the convergence rate important?





# **Applied Numerical Optimization**

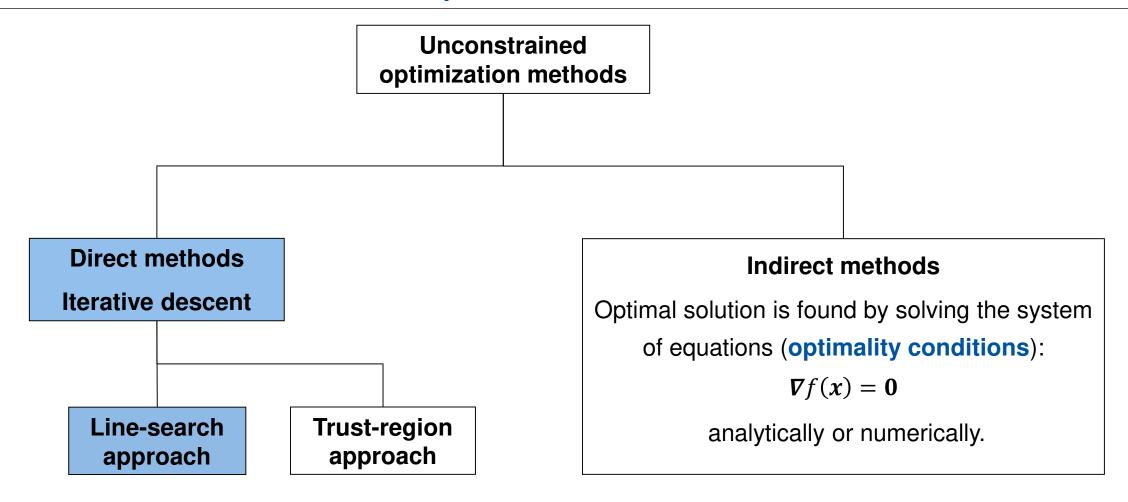
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Line search: basic idea and step length





# **Solution Methods for Unconstrained Optimization**





# **Direct Methods – Line-Search Approach**

#### Definition (descent direction):

A vector p is called descent direction at  $x^{(k)}$ , if  $\nabla f(x^{(k)})^T p < 0$  holds.

#### Basic algorithm (line-search):

- 1. Choose a descent direction,  $p^{(k)}$ , such that  $\nabla f(x^{(k)})^T p^{(k)} < 0$
- 2. Determine a step length  $\alpha_k$
- 3. Set  $x^{(k+1)} = x^{(k)} + \alpha_k p^{(k)}$

# $\|\alpha_{k} \boldsymbol{p}^{(k)}\|$ $\boldsymbol{x}^{(k)}$ $\boldsymbol{x}^{(k)}$ $\Omega = R^{n}$

#### Open issues:

- Determination of the descent direction  $p^{(k)}$ ?
- Calculation of the step length  $\alpha_k$ ?





# Calculation of Step Length $\alpha_k$

#### The exact line search strategy:

1. Define the one-dimensional function along the descent direction  $p^{(k)}$ .

$$\phi(\alpha) = f(\mathbf{x}^{(k)} + \alpha \mathbf{p}^{(k)})$$

2. Solve the one-dimensional minimization problem

$$\min_{\alpha>0}\phi(\alpha)$$

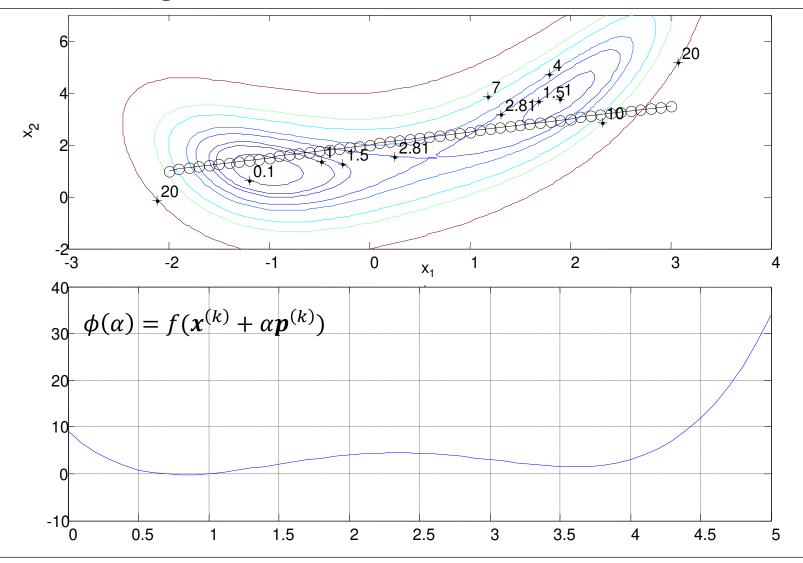
#### Remarks

- 1. Naively speaking it would be ideal to globally minimize  $\phi(\alpha)$ . Generally, it is very expensive to find this solution. It is not necessarily a good idea since the search is one-dimensional
- 2. One could also search for some local solution. But this is often also too expensive (need function and/or gradient evaluations at a number of points).
- 3. Practical strategies (so-called non-exact LS): find  $\alpha$  such that  $f(x^{(k+1)})$  becomes as small as possible with minimal effort.





# **Practical Line-Search Strategies**







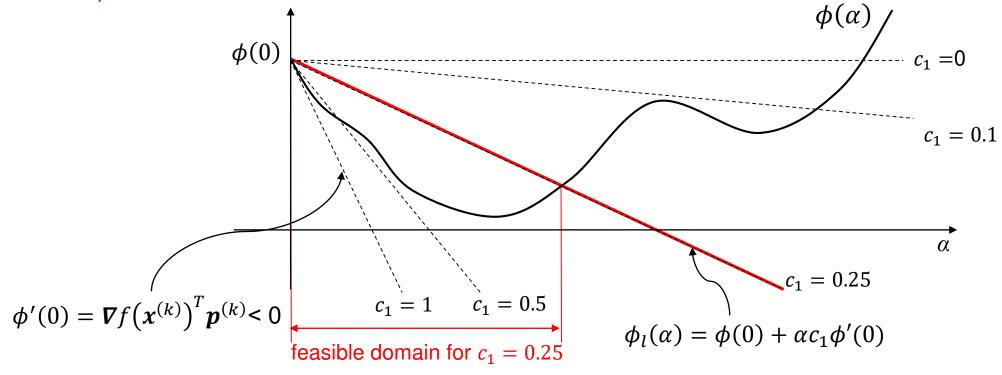
## **Armijo Condition**

#### Theorem<sup>[1]</sup>:

Let f be continuously differentiable,  $\mathbf{p}^{(k)}$  a descent direction, and let  $c_1 \in (0,1)$  be given. Then there exists an  $\alpha > 0$ , such that for  $\phi(\alpha) \coloneqq f(\mathbf{x}^{(k)} + \alpha \mathbf{p}^{(k)})$ , the condition  $\phi(\alpha) \le \phi(0) + \alpha c_1 \phi'(0)$  holds.

#### Geometrical interpretation:

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# **Simple Line-Search Algorithm**

#### Remarks:

1. The choice of a step length, which fulfills the Armijo condition guarantees the descent of f:

$$\phi'(0) = \nabla f(x^{(k)})^T p^{(k)} < 0 \ (p^{(k)} \text{ is a descent direction})$$

$$\phi(\alpha) \le \phi(0) + \alpha c_1 \phi'(0) \qquad \Rightarrow \phi(\alpha) < \phi(0)$$

$$\Rightarrow \text{ a descent is guaranteed!}$$

- 2. The choice of  $c_1$  is crucial:
  - Large  $c_1$  leads to small values of  $\alpha$ , such that  $x^{(k+1)} \approx x^{(k)}$ .
  - Small  $c_1$  potentially results in small reduction of f and therefore slower convergence

```
Simple line-search algorithm: choose \alpha_1 > 0; \rho, c_1 \in (0,1) set \alpha = \alpha_1 repeat \alpha \leftarrow \rho \alpha until \phi(\alpha) \leq \phi(0) + \alpha c_1 \phi'(0)
```





# **Improved Line-Search Algorithm**

choose  $\alpha_0 > 0$  and  $c_1 \in (0,1)$ 

if 
$$\phi(\alpha_0) \leq \phi(0) + \alpha_0 c_1 \phi'(0)$$
 STOP, else

find a better  $\alpha \in (0, \alpha_0)$  through *quadratic interpolation* of available data:

$$\alpha_1 = -\frac{\phi'(0)\alpha_0^2}{2[\phi(\alpha_0) - \phi(0) - \phi'(0)\alpha_0]}$$

if 
$$\phi(\alpha_1) \leq \phi(0) + \alpha_1 c_1 \phi'(0)$$
 STOP, else

find a better  $\alpha \in (0, \alpha_1)$  through *cubic interpolation* of available data (how ?)

repeat the procedure of *cubic interpolation*, until the condition is fulfilled





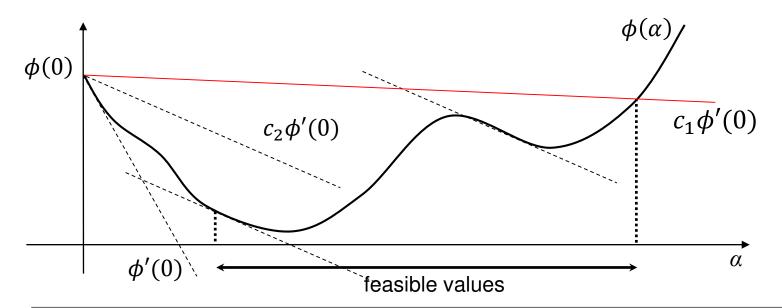
#### **Wolfe Conditions**

#### Theorem<sup>[1]</sup>:

Let f be continuously differentiable,  $p^{(k)}$  a descent direction and  $c_1 \in (0,1)$ ,  $c_2 \in (c_1,1)$ . Then, there exists an  $\alpha > 0$ , such that  $\phi(\alpha) \leq \phi(0) + \alpha c_1 \phi'(0)$ 

$$\phi'(\alpha) \ge c_2 \phi'(0)$$
 (slope condition)

**Geometric interpretation:** → guarantee minimum step length!



#### Relevance:

Wolfe Conditions promote convergence to a stationary point<sup>[1]</sup>





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#### **Check Yourself**

- Explain the basic ideas of the line-search method.
- What is a descent direction? How it is defined?
- Explain the Armijo-rule and its potential drawbacks?
- Explain the Wolfe conditions and the advantage compared to Armijo's rule.







# **Applied Numerical Optimization**

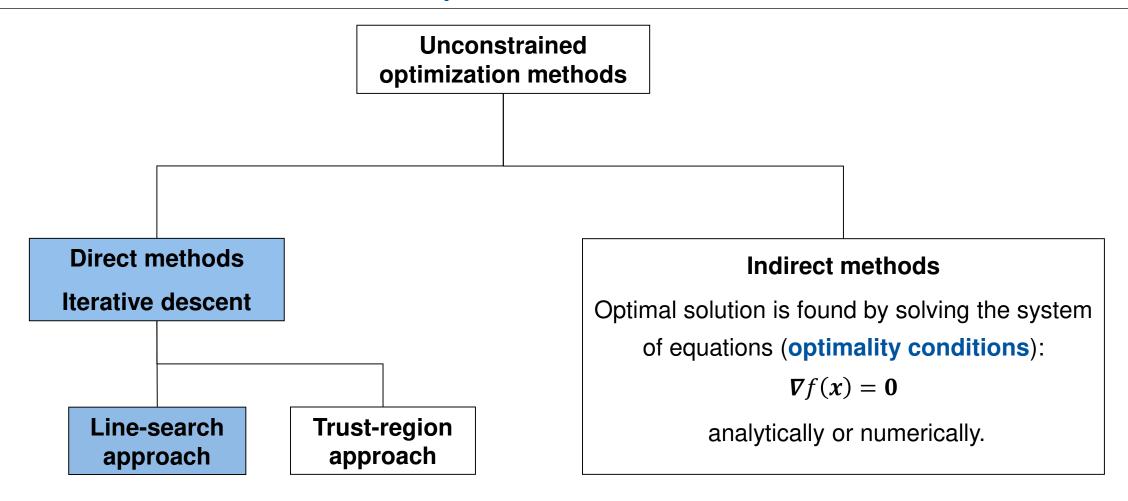
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Line search: simple directions





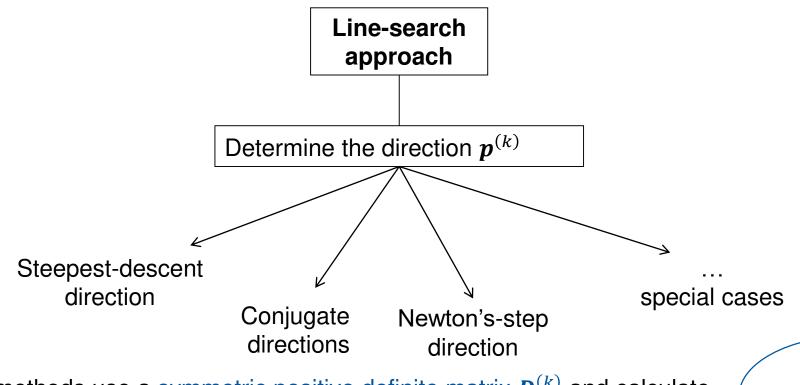
# **Solution Methods for Unconstrained Optimization**





#### **Determination of a Descent Direction: A Toolbox**

Line-search approaches differ from each other with respect to the determination of descent direction and step length.



Many gradient methods use a symmetric positive definite matrix  $\mathbf{D}^{(k)}$  and calculate

 $\boldsymbol{x}^{(k+1)} = \boldsymbol{x}^{(k)} - \alpha_k \boldsymbol{D}^{(k)} \boldsymbol{\nabla} f(\boldsymbol{x}^{(k)})$ 

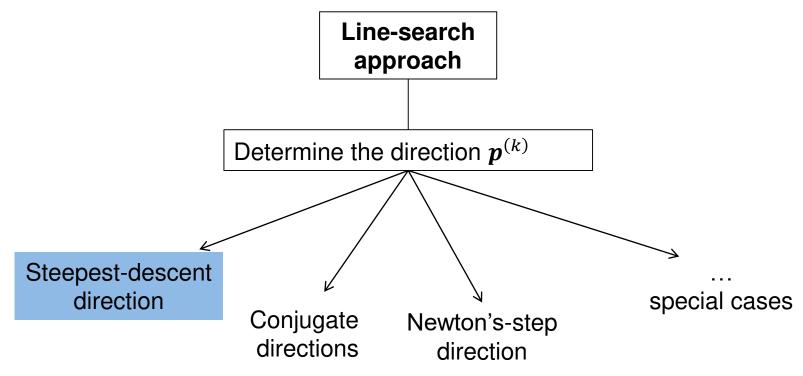
Extra work: prove that it guarantees descent!





#### **Determination of a Descent Direction: A Toolbox**

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$$\boldsymbol{x}^{(k+1)} = \boldsymbol{x}^{(k)} - \alpha_k \boldsymbol{D}^{(k)} \boldsymbol{\nabla} f(\boldsymbol{x}^{(k)})$$





# **Steepest-Descent Direction (1)**

Taylor series: 
$$f(\mathbf{x}^{(k)} + \alpha \mathbf{p}^{(k)}) = f(\mathbf{x}^{(k)}) + \left[\alpha \nabla f(\mathbf{x}^{(k)})^T \mathbf{p}^{(k)}\right] + O(\alpha^2)$$

The rate of change of f at  $x^{(k)}$  along the direction  $p^{(k)}$  is the coefficient in the linear term:

$$\nabla f(x^{(k)})^T p^{(k)}$$

The unit direction  $p^{(k)}$  with the **highest rate of change** is the solution of the following problem

$$\min_{\boldsymbol{p}^{(k)} \in R^n} \nabla f(\boldsymbol{x}^{(k)})^T \boldsymbol{p}^{(k)} \quad \text{s. t. } \|\boldsymbol{p}^{(k)}\| = 1$$

Note that  $\nabla f(x^{(k)})^T p^{(k)} = ||\nabla f(x^{(k)})|| ||p^{(k)}|| \cos(\theta)$ 

The solution of the problem is achieved for  $cos(\theta) = -1 \Rightarrow \theta = \pi$ 

$$\Rightarrow \boldsymbol{p}^{(k)} = -\boldsymbol{\nabla} f(\boldsymbol{x}^{(k)}) / \|\boldsymbol{\nabla} f(\boldsymbol{x}^{(k)})\|$$

The choice of  $\mathbf{D}^{(k)}$  is the identity matrix  $\mathbf{I}$ .





# **Steepest-Descent Direction (2)**

descent direction:  $\nabla f(x^{(k)})^T p^{(k)} < 0$  $f(\mathbf{x}^{(k)}) > C$  $f(\mathbf{x}^{(k)}) = C$  $f(\mathbf{x}^{(k)}) < C$  $\nabla f(x^{(k)})^T$  $\nabla f(x^{(k)})$  $\Theta_1$  $\nabla f(\mathbf{x}^{(k)})^T \mathbf{p}^{(k)} = \|\nabla f(\mathbf{x}^{(k)})\| \|\mathbf{p}^{(k)}\| \cos(\theta)$  $\Theta_2$  $\nabla f(x^{(k)})^{T} p^{(2)}$  $x^{(k)}$  $p^{(2)}$  $-\nabla f(x^{(k)})$ 





# **Method of Steepest-Descent**

#### Algorithm:

choose  $x^{(0)}$ 

for k=0,1,...

if 
$$\|\nabla f(x^{(k)})\| \le \varepsilon$$
 stop, else

$$\mathbf{set}\; \boldsymbol{p}^{(k)} = -\boldsymbol{\nabla} f(\boldsymbol{x}^{(k)})$$

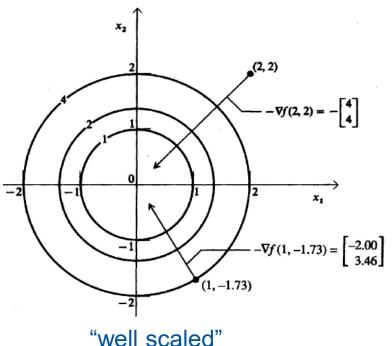
determine the step length  $\alpha_k$  (e.g. using the Armijo rule)

set 
$$x^{(k+1)} = x^{(k)} + \alpha_k p^{(k)}$$

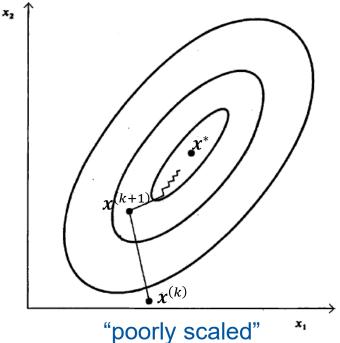
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#### end for







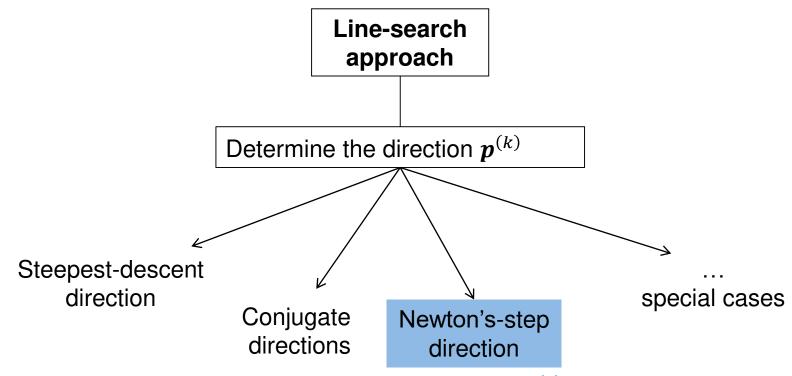
Directions become perpendicular





#### **Determination of a Descent Direction: A Toolbox**

Line-search approaches differ from each other with respect to the determination of descent direction and step length.



Many gradient methods use a symmetric positive definite matrix  $\mathbf{D}^{(k)}$  and calculate

$$\boldsymbol{x}^{(k+1)} = \boldsymbol{x}^{(k)} - \alpha_k \boldsymbol{D}^{(k)} \boldsymbol{\nabla} f(\boldsymbol{x}^{(k)})$$





#### **Newton's Descent Direction**

### Quadratic approximation of f at $x^{(k+1)}$

$$m(\boldsymbol{x}^{(k+1)}) = f(\boldsymbol{x}^{(k)}) + \nabla f(\boldsymbol{x}^{(k)})^T (\boldsymbol{x}^{(k+1)} - \boldsymbol{x}^{(k)}) + \frac{1}{2} (\boldsymbol{x}^{(k+1)} - \boldsymbol{x}^{(k)})^T \nabla^2 f(\boldsymbol{x}^{(k)}) (\boldsymbol{x}^{(k+1)} - \boldsymbol{x}^{(k)})$$

$$(1st nec. opt. cond. for m)$$

$$0 = \nabla m(\boldsymbol{x}^{(k+1)}) = \nabla f(\boldsymbol{x}^{(k)}) + \nabla^2 f(\boldsymbol{x}^{(k)}) (\boldsymbol{x}^{(k+1)} - \boldsymbol{x}^{(k)})$$

$$\Rightarrow \boldsymbol{x}^{(k+1)} = \boldsymbol{x}^{(k)} - [\nabla^2 f(\boldsymbol{x}^{(k)})]^{-1} \nabla f(\boldsymbol{x}^{(k)})$$

$$\Rightarrow \boldsymbol{p}^{(k)} = -[\nabla^2 f(\boldsymbol{x}^{(k)})]^{-1} \nabla f(\boldsymbol{x}^{(k)})$$
minimum of the quadratic approximation quadratic

The choice of  $\mathbf{D}^{(k)}$  is the inverse of the Hessian

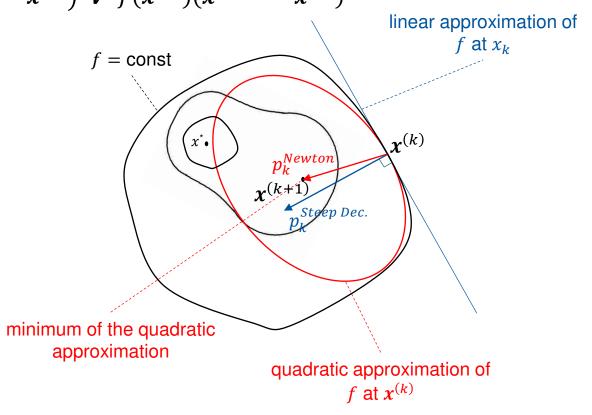


Fig.: Comparison of steepest-descent with Newton's method from viewpoint of objective function approximation





#### **Newton's Method**

#### Algorithm:

choose 
$$x^{(0)}$$
 for  $k=0,1,...$  if  $\|\nabla f(x^{(k)})\| \leq \varepsilon$  stop, else 
$$\operatorname{set} \boldsymbol{p}^{(k)} = - \left[\nabla^2 f(x^{(k)})\right]^{-1} \nabla f(x^{(k)})$$
 
$$\operatorname{set} x^{(k+1)} = x^{(k)} + \boldsymbol{p}^{(k)}$$

#### end for

#### Remarks:

1. line-search?

$$\mathbf{x}^{(k+1)} = \mathbf{x}^{(k)} + \alpha_k \mathbf{p}^{(k)}$$
$$\mathbf{p}^{(k)} = -[\nabla^2 f(\mathbf{x}^{(k)})]^{-1} \nabla f(\mathbf{x}^{(k)})$$
$$\alpha_k = 1$$

- 2. (+) locally quadratic convergence, if  $x^{(k)}$  close to  $x^*$  (-)  $2^{nd}$  derivatives & inversion (expensive for large system of equations)
- 3. If *f* is quadratic, the algorithm converges in one iteration.
- 4. Convergence to a minimum is not guaranteed! Why?





#### **Check Yourself**

- Explain the basic ideas of the line-search method.
- Explain the steepest descent method.
- What additional requirements puts Newton's method on the objective function?
- Explain the Newton direction. Is it better than other descent directions? Why is the Newton step-length equal to one?





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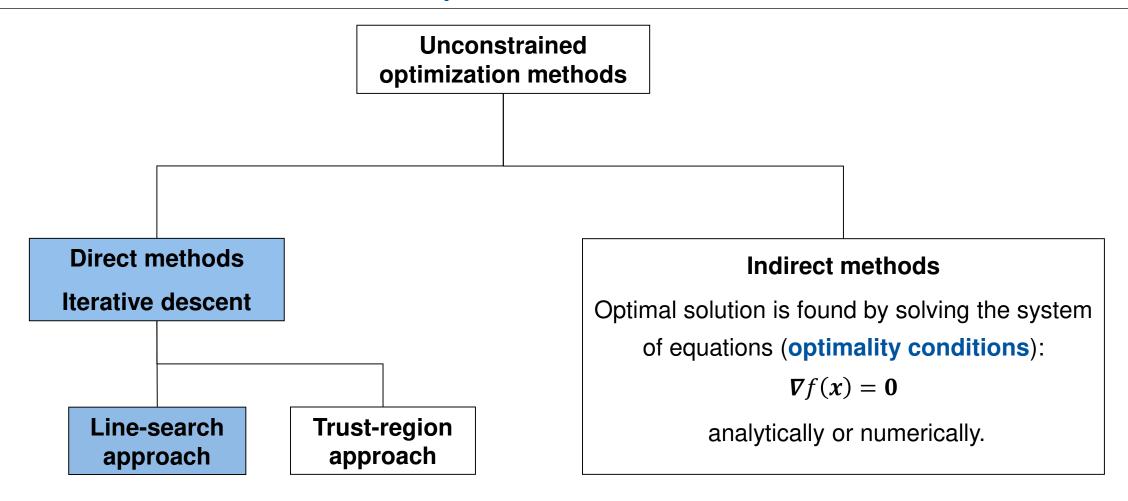
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Line search: complexity and examples





# **Solution Methods for Unconstrained Optimization**





# **Complexity Analysis**

Nesterov (2004) proves: "In general, optimization problems are unsolvable" \*

Let F denote a class of problems, e.g., Lipschitz-continuous functions with Lipschitz-constant L, i.e., |f(x) - f(y)| < L||x - y||, L is assumed to be fixed for all  $P \in F$ .

"Performance of a method M on a problem  $P \in F$  is the total amount of computational effort that is required by M to solve P." \*

"To solve the problem means to find an approximate solution to P with an accuracy  $\varepsilon > 0$ ." \*

For unconstrained problems, the accuracy  $\varepsilon > 0$  can be defined as the **norm of the objective's gradient**.

\* Yurii Nesterov, Introductory Lectures on Convex Optimization – A Basic Course, Kluwer Academic Publishers, (2004)





# **Complexity Analysis – Measuring Computational Effort**

#### Unit of measurement: Query to an oracle

It is assumed that the objective function is unknown and that the algorithm solves the optimization problem by *querying an oracle* for local information about the unknown objective function. An oracle is simply a "black box" capable of answering any query of the form:

Given x return the value f(x) (Zeroth-order oracle)

• Given x return f(x) and gradient  $\nabla f(x)$  (First-order oracle)

Given x return f(x),  $\nabla f(x)$  and Hessian  $\nabla^2 f(x)$  (Second-order oracle)

Analytical Complexity: The smallest number of queries to an oracle to solve Problem P to accuracy  $\varepsilon$ . [1]

**Arithmetical Complexity**: The smallest number of arithmetic operations (including work of the oracle and work of method), required to solve problem P up to accuracy  $\varepsilon$ . [1]





"Oracle"

# **Analytical Complexity of Steepest Descent Method**

#### Algorithm:

```
choose x^{(0)} for k=0,1,... if \|\nabla f(x^{(k)})\| \leq \varepsilon stop, else set p^{(k)} = -\nabla f(x^{(k)}) determine the step length \alpha_k (e.g. using the Armijo rule) set x^{(k+1)} = x^{(k)} + \alpha_k p^{(k)}
```

#### end for

- Problem class: f is continuously differentiable and  $\nabla f(x)$  is Lipschitz-continuous with fixed Lipschitz constant L, i.e.,  $\|\nabla f(x) \nabla f(y)\| < L\|x y\|$
- First-order oracle: returns f(x) and gradient  $\nabla f(x)$
- Worst-case analytical complexity (queries to oracle):  $O\left(\frac{1}{\varepsilon^2}\right)$

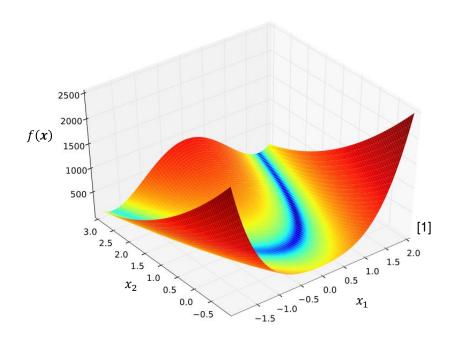




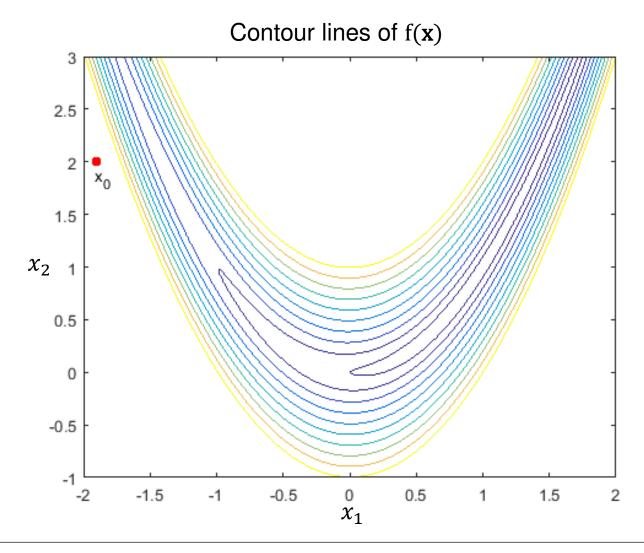
## **Rosenbrock Function**

$$\min_{\mathbf{x} \in R^2} f(\mathbf{x}) = 100(x_2 - x_1^2)^2 + (1 - x_1)^2$$

• Solution point is  $x = (1,1)^T$  - why?



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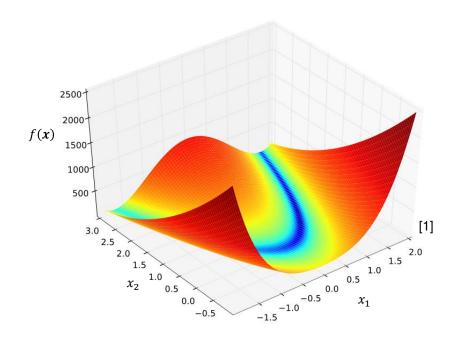




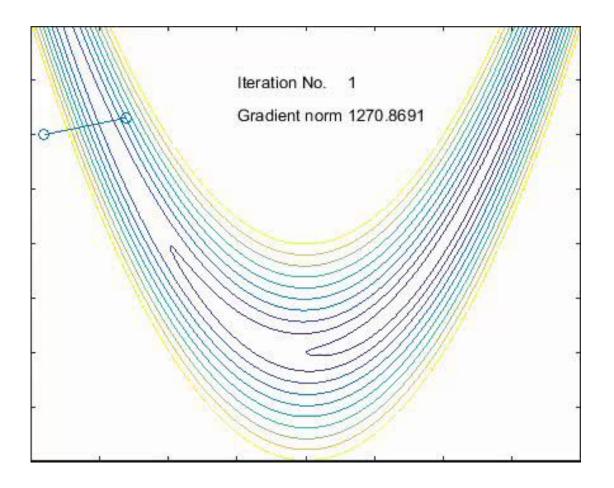
# **Illustration of Convergence (1)**

$$\min_{\mathbf{x} \in R^2} f(\mathbf{x}) = 100(x_2 - x_1^2)^2 + (1 - x_1)^2$$

Steepest descent with Armijo line-search



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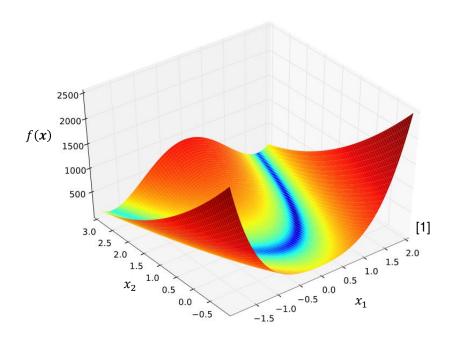


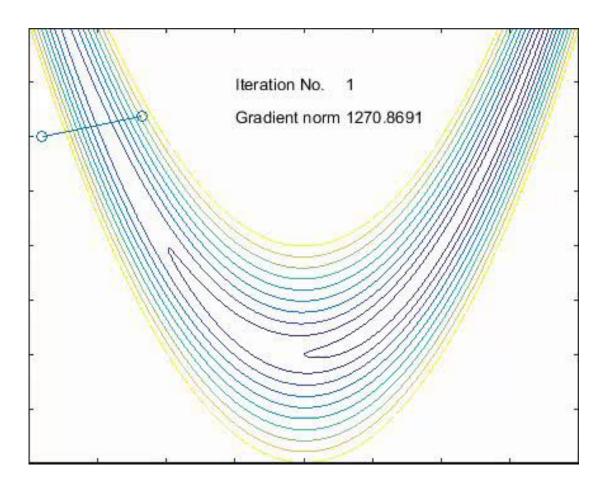


# **Illustration of Convergence (2)**

$$\min_{\mathbf{x} \in R^2} f(\mathbf{x}) = 100(x_2 - x_1^2)^2 + (1 - x_1)^2$$

Steepest descent with Wolfe line-search









## **Analytical Complexity of Newton's Method**

- Problem class: f is twice continuously differentiable and  $\nabla^2 f(x)$  is Lipschitz-continuous with fixed Lipschitz constant L, i.e.,  $\|\nabla^2 f(x) \nabla^2 f(y)\| < L\|x y\|$
- Second-order oracle: returns f(x),  $\nabla f(x)$  and Hessian  $\nabla^2 f(x)$
- Quadratic approximation of f around  $x^{(k)}$ , line search

$$m(x^{(k+1)}) = f(x^{(k)}) + \nabla f(x^{(k)})^{T} (x^{(k+1)} - x^{(k)}) + \frac{1}{2} (x^{(k+1)} - x^{(k)})^{T} \nabla^{2} f(x^{(k)}) (x^{(k+1)} - x^{(k)})$$

- Worst-case analytical complexity:  $O\left(\frac{1}{\varepsilon^{2-\tau}}\right)$ ,  $1 > \tau > 0$ , arbitrary but fixed for a given problem
- Quadratic approximation of f around  $x^{(k)}$  with cubic regularization, line search

$$m_{regularized}(x^{(k+1)}) = m(x^{(k+1)}) + \frac{1}{3}\sigma_k ||x^{(k+1)} - x^{(k)}||^3$$

• Worst-case analytical complexity:  $O\left(\frac{1}{\varepsilon^{3/2}}\right)$ 

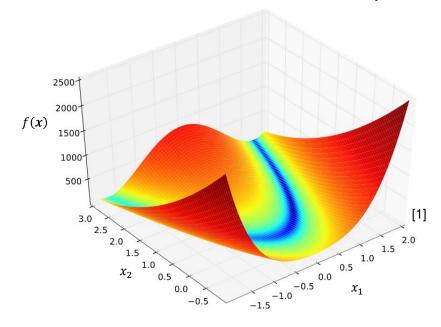


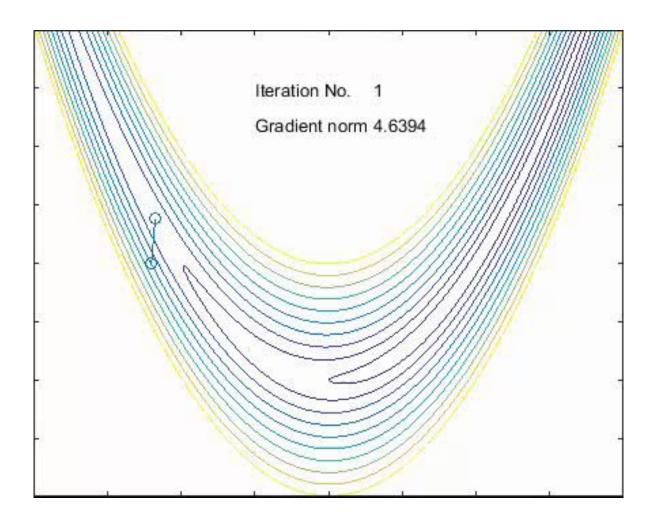


# **Illustration of Convergence (3)**

$$\min_{\mathbf{x} \in R^2} f(\mathbf{x}) = 100(x_2 - x_1^2)^2 + (1 - x_1)^2$$

**Modified Newton** (Armijo line-search; if Hessian is < 0 switch to steepest descent)









## **Check Yourself**

- What does the term complexity analysis refer to?
- What is the difference of analytical and arithmetic complexity
- Which method has better analytical complexity: Newton vs. steepest descent?







# **Applied Numerical Optimization**

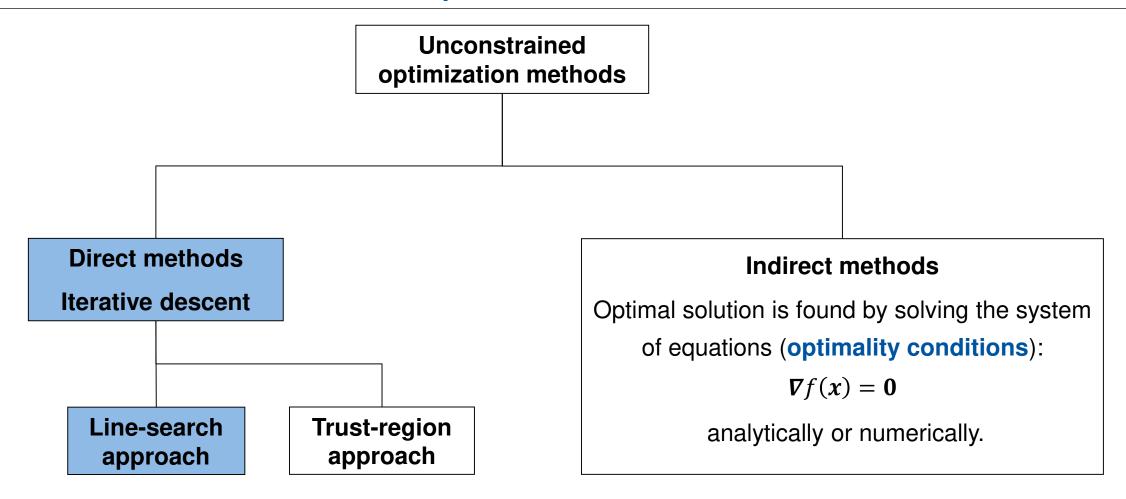
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Line search: advanced directions





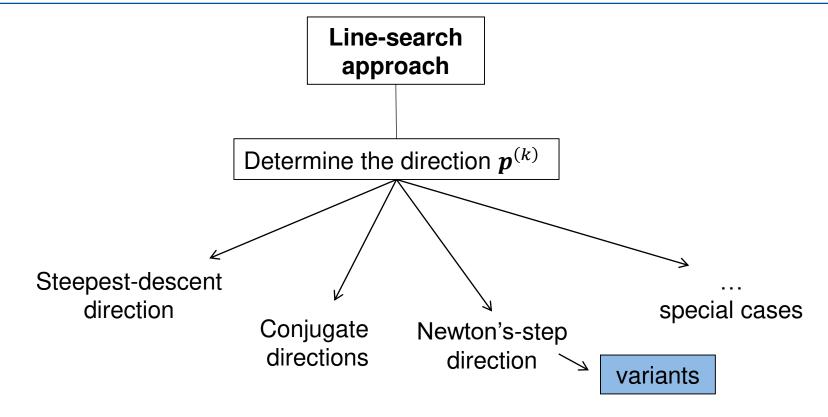
# **Solution Methods for Unconstrained Optimization**





### **Determination of a Descent Direction: A Toolbox**

Line-search approaches differ from each other with respect to the determination of descent direction and step length.







## **Inexact Newton Method (1)**

**Define:** 
$$f^{(k)} = f(x^{(k)})$$
 and  $g^{(k)} := \nabla f(x^{(k)})$  and  $H^{(k)} := \nabla^2 f(x^{(k)})$ 

#### From Newton's method:

$$\boldsymbol{x}^{(k+1)} = \boldsymbol{x}^{(k)} + \boldsymbol{p}^{(k)}$$
$$\left[\boldsymbol{\nabla}^{2} f(\boldsymbol{x}^{(k)})\right] \boldsymbol{p}^{(k)} = -\boldsymbol{\nabla} f(\boldsymbol{x}^{(k)}) \Rightarrow \boldsymbol{H}^{(k)} \boldsymbol{p}^{(k)} = -\boldsymbol{g}^{(k)}$$

#### Idea:

• The linear equation system,  $H^{(k)}p^{(k)} = -g^{(k)}$ , is solved approximately by an iterative method, e.g., by CG (conjugate gradients) if  $H^{(k)}$  is positive definite.

#### **Comments:**

- LU- or Cholesky-decomposition very high computational effort!
- Large errors occur for ill-conditioned problems.
- The exact solution is not needed.





## **Inexact Newton Method (2)**

### Newton-CG method:

Newton's method

CG method to determine  $oldsymbol{p}^{(k)}$  approximately

### Algorithm:

choose  $x^{(0)}$ 

for k=0,1,...

end for

if  $\|\nabla f(x^{(k)})\| \le \varepsilon$  stop, else

calculate  $g^{(k)} \coloneqq \nabla f(x^{(k)})$  and  $H^{(k)} \coloneqq \nabla^2 f(x^{(k)})$ 

solve  $H^{(k)}p^{(k)} = -g^{(k)}$  for  $p^{(k)}$  with CG method

 $\mathbf{set} \ \boldsymbol{x}^{(k+1)} = \boldsymbol{x}^{(k)} + \alpha_k \boldsymbol{p}^{(k)}$ 

Some line search strategy is needed. (Why?)





## **Modified Newton Method**

### **Motivation:** What if

- $H^{(k)}$  is singular or almost singular (poorly conditioned)?
- $H^{(k)}$  is not positive definite?

$\mathbf{H}^{(k)} \coloneqq \mathbf{\nabla}^2 f(\mathbf{x}^{(k)})$	$\mathbf{H}^{(k)}\coloneqq$	$\nabla^2 f(x^{(k)})$
--	-----------------------------	-----------------------

Idea	Approximations
replace $ extbf{ extit{H}}^{(k)}$ by the approximation $ extbf{ extit{B}}^{(k)} pprox  extbf{ extit{H}}^{(k)}$	$m{B}^{(k)} = m{H}^{(k)} + m{E}^{(k)}$ with $m{E}^{(k)} =  au_k m{I}$ , $ au_k \geq 0$ smartly chosen
$\boldsymbol{B}^{(k)}\boldsymbol{p}^{(k)} = -\boldsymbol{g}^{(k)}$	converges to steepest descent for $\tau_k \to \infty$
$\pmb{x}^{(k+1)} = \pmb{x}^{(k)} + \alpha_k \pmb{p}^{(k)}$ , ( $\alpha_k$ from the line-search)	

Alternatives exist, e.g., see [1]





# **Quasi-Newton Methods (1)**

# **Idea**: Reduce complexity by simplified calculation of $H^{(k)}$ (Davidon):

- replace  $H^{(k)}$  by an approximation  $B^{(k)}$ .
- instead of calculating  $B^{(k)}$ , we look for a simple update using information from the last iterations.

$$\boldsymbol{H}^{(k)} \coloneqq \boldsymbol{\nabla}^2 f(\boldsymbol{x}^{(k)})$$

$$\boldsymbol{g}^{(k)} \coloneqq \boldsymbol{\nabla} f(\boldsymbol{x}^{(k)})$$

$$f^{(k)} \coloneqq f(\mathbf{x}^{(k)})$$

## Approach:

- Consider quadratic approximation of f at  $\mathbf{x}^{(k)}$ ,  $m^{(k)}(\mathbf{p}) = f^{(k)} + \mathbf{g}^{(k)}\mathbf{p} + \frac{1}{2}\mathbf{p}^T\mathbf{g}^{(k)}\mathbf{p}$ .
- First order optimality condition:  $\mathbf{p}^{(k)} = -\mathbf{B}^{(k)^{-1}}\mathbf{g}^{(k)}$
- By convexity necessary and sufficient for minimization of  $m^{(k)}(\mathbf{p})$ .

symmetric positive definite

• Construct the quadratic approximation at  $x^{(k+1)} = x^{(k)} + \alpha_k p^{(k)}$ ,

$$m^{(k+1)}(\mathbf{p}) = f^{(k+1)} + \mathbf{g}^{(k+1)^T} \mathbf{p} + \frac{1}{2} \mathbf{p}^T \mathbf{B}^{(k+1)} \mathbf{p}$$

• What conditions must  $B^{(k+1)}$  satisfy?



## **Quasi-Newton Methods (2)**

## Conditions on $B^{(k+1)}$ :

Gradient of  $m^{(k+1)}$  at  $x^{(k)}$  and  $x^{(k+1)}$  must be equal to gradient of f.

$\nabla m^{(k+1)}(p) = g^{(k+1)} + B^{(k+1)}p$	
At $x = x^{(k+1)}$ , $p = 0$	At $\boldsymbol{x} = \boldsymbol{x}^{(k)}$ , $\boldsymbol{p} = -\alpha_k \boldsymbol{p}^{(k)}$
We want $\nabla m^{(k+1)}(0) = g^{(k+1)}$	We want $\nabla m^{(k+1)} \left( -\alpha_k \boldsymbol{p}^{(k)} \right) = \boldsymbol{g}^{(k)}$
	$\Rightarrow \boldsymbol{g}^{(k+1)} - \alpha_k \boldsymbol{B}^{(k+1)} \boldsymbol{p}^{(k)} = \boldsymbol{g}^{(k)}$
Automatically satisfied	$\Rightarrow \mathbf{B}^{(k+1)} \underbrace{\alpha_k \mathbf{p}^{(k)}}_{=\mathbf{x}^{(k+1)} - \mathbf{x}^{(k)}} = \mathbf{g}^{(k+1)} - \mathbf{g}^{(k)}$
	$\Rightarrow \boxed{ m{B}^{(k+1)} m{s}^{(k)} = m{y}^{(k)} },  ext{ where } m{s}^{(k)} = m{x}^{(k+1)} - m{x}^{(k)}  ext{ and } m{y}^{(k)} = m{g}^{(k+1)} - m{g}^{(k)} $

Since,  $\mathbf{B}^{(k+1)}$  is symmetric positive definite:  $\mathbf{s}^{(k)^T}\mathbf{B}^{(k+1)}\mathbf{s}^{(k)} > 0$ ,  $\forall \mathbf{s}^{(k)} \neq 0 \Rightarrow \mathbf{s}^{(k)^T}\mathbf{y}^{(k)} > 0$ 



 $\star$  Wolfe conditions (line-search) guarantee these constraints for all f, even when f is non-convex.





## **Quasi-Newton Methods (3)**

## Conditions on $B^{(k+1)}$ :

 $\mathbf{B}^{(k+1)}\mathbf{s}^{(k)} = \mathbf{y}^{(k)}$  gives many solutions for  $\mathbf{B}^{(k+1)}$ 

• Unique solution:  $B^{(k+1)}$  should be close to  $B^{(k)}$ 

$$\min_{\boldsymbol{B}} \left\| \boldsymbol{B} - \boldsymbol{B}^{(k)} \right\|_{W} \leftarrow \text{weighted Frobenius-Norm}$$
s. t.  $\boldsymbol{B}^{T} = \boldsymbol{B}$ 

$$\left\| \boldsymbol{A} \right\|_{W} = \left\| W^{1/2} \boldsymbol{A} W^{1/2} \right\|_{F}, \text{ for any } W \text{ s.t. } W y_{k} = s_{k}$$

$$\left\| \boldsymbol{B} \boldsymbol{s}^{(k)} = \boldsymbol{y}^{(k)} \right\|_{F}^{2} : R^{n \times n} \rightarrow R_{\geq 0}, \left\| \boldsymbol{C} \right\|_{F}^{2} = \sum_{i=1}^{n} \sum_{j=1}^{n} c_{ij}^{2}$$

$$\Rightarrow B^{(k+1)} = \left(I - \frac{1}{y^{(k)^T} s^{(k)}} y^{(k)} s^{(k)^T}\right) B^{(k)} \left(I - \frac{1}{y^{(k)^T} s^{(k)}} s^{(k)} y^{(k)^T}\right) + \frac{1}{y^{(k)^T} s^{(k)}} y^{(k)} y^{(k)^T} \quad \to \text{DFP formula}$$

$$\Rightarrow B^{(k+1)^{-1}} = \left(I - \frac{1}{y^{(k)^T} s^{(k)}} s^{(k)} y^{(k)^T}\right) B^{(k)^{-1}} \left(I - \frac{1}{y^{(k)^T} s^{(k)}} y^{(k)} s^{(k)^T}\right) + \frac{1}{y^{(k)^T} s^{(k)}} s^{(k)} s^{(k)^T} \rightarrow \text{BFGS formula}$$





### **Check Yourself**

- Explain the inexact Newton method.
- What is the main idea of the modified and quasi-Newton methods? Why are these methods advantageous?
- Why is it necessary to introduce a step-length control mechanism (line-search) into modified and quasi-Newton methods?







# **Applied Numerical Optimization**

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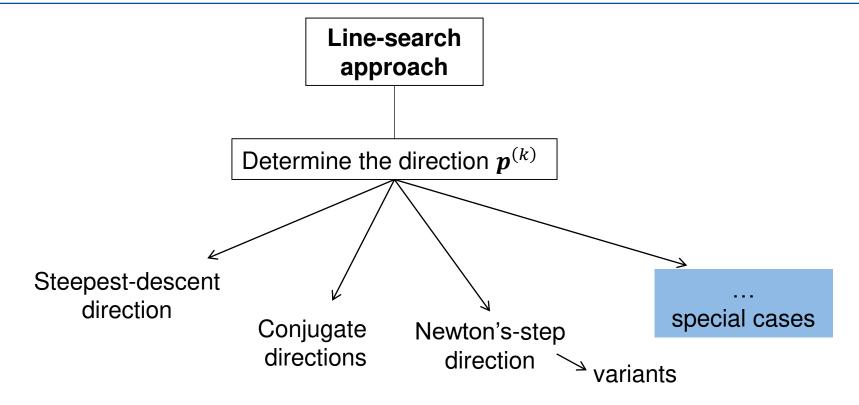
Parameter estimation





### **Determination of a Descent Direction: A Toolbox**

Line-search approaches differ from each other with respect to the determination of descent direction and step length.







# **Regression Problems: Least-Squares Formulation**

### Example:

Consider a batch reactor with the reaction  $A \to B$  at constant temperature  $T_R$ . The reagent concentration  $c_A$  is measured at time instants  $t_j$ .

The reaction is of first order, therefore we can write the analytic solution:

$$\left. \frac{dc_A}{dt} \right|_t = -k \cdot c_A(t) \quad \rightarrow \quad c_A(t) = c_A|_{t=0} \cdot e^{-kt}$$

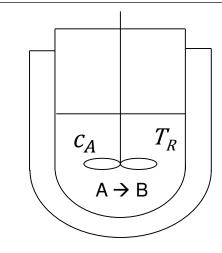
The reaction constant k and the reagent concentration at initial time  $c_A(t=0)$  are unknown and should be determined from the measurements.

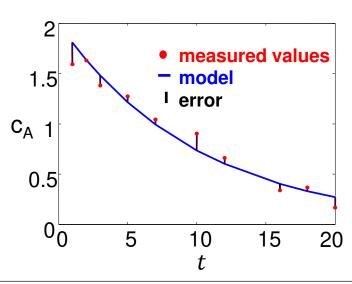
Optimization formulation uses  $x_1 = c_A(t = 0), x_2 = k$ :

$$c_{A, \, \mathrm{theoretical}}(t_j) = \varphi(x, t_j) = x_1 \cdot e^{x_2 \cdot t_j}$$
 model  $c_{A, \, \mathrm{measured}}(t_i) = y_i$  measurement

$$\varepsilon_i = y_i - \varphi(x, t_i), \forall j = 1, ..., m$$
 residual (error)

$$\min_{\boldsymbol{x} \in R^2} \frac{1}{2} \|\boldsymbol{\varepsilon}(\boldsymbol{x})\|_2^2 = \frac{1}{2} \Sigma_j \left( y_j - \phi(\boldsymbol{x}, t_j) \right)^2$$









### **Gauss-Newton Method**

• 
$$\min_{\mathbf{x} \in R^2} f(\mathbf{x}) = \min_{\mathbf{x} \in R^2} \frac{1}{2} \| \mathbf{\varepsilon}(\mathbf{x}) \|_2^2 = \min_{\mathbf{x} \in R^2} \frac{1}{2} \mathbf{\varepsilon}(\mathbf{x})^T \mathbf{\varepsilon}(\mathbf{x})$$

• Define:  $I(x) := \nabla \varepsilon(x) \in \mathbb{R}^{m \times 2}$ 

$$\Rightarrow \nabla f(x) = J(x)^T \varepsilon(x)$$

$$\Rightarrow \nabla^2 f(x) = J(x)^T J(x) + \sum_{j=1}^m \varepsilon_j(x) \nabla^2 \varepsilon_j(x)$$

- The Hessian can be approximated by the first term in case of almost linear problems (i.e.,  $\nabla^2 \varepsilon_j(x) = 0$ ) or good starting values (i.e., small  $\varepsilon_j(x)$ )
- Newton's direction:  $\nabla^2 f(x^{(k)}) p_k = -\nabla f(x^{(k)})$
- With Hessian approximation:  $J^{(k)}^T J^{(k)} p^{(k)} = -J^{(k)}^T \varepsilon^{(k)}$





### **Remarks on Gauss-Newton Method**

• If  $J^{(k)}$  has full-rank,  $p^{(k)}$  is always a descent direction

$$p^{(k)^{T}} \cdot \nabla f(x^{(k)}) = p^{(k)^{T}} \cdot J^{(k)^{T}} \varepsilon^{(k)} = -p^{(k)^{T}} \cdot J^{(k)^{T}} \cdot J^{(k)} \cdot p^{(k)} = -\|J^{(k)} \cdot p^{(k)}\|_{2}^{2} < 0,$$

The inequality is strict unless  $\boldsymbol{J}^{(k)} \cdot \boldsymbol{p}^{(k)} = 0 \Leftrightarrow \boldsymbol{J}^{(k)} \boldsymbol{\varepsilon}^{(k)} = \boldsymbol{\nabla} f_k = 0$ .  $\boldsymbol{\leftarrow}$  Optimum

- In descent-direction  $m{p}^{(k)}$ , the step-length is determined as per the Wolfe-conditions
- For linear models the Jacobian *I* matrix is constant.
- The condition of minimization corresponds to the normal equations
- If  $J^{(k)}$  is singular or almost singular, the descent direction  $p^{(k)}$  is, usually, not reliable. The method converges very poorly. Quasi-Newton methods are therefore more efficient.
- It is a local method





## **Check Yourself**

- Where can least-squares problems be applied?
- When is a least-squares problem linear or nonlinear?
- What is the key idea of the Gauss-Newton method?





# **Applied Numerical Optimization**

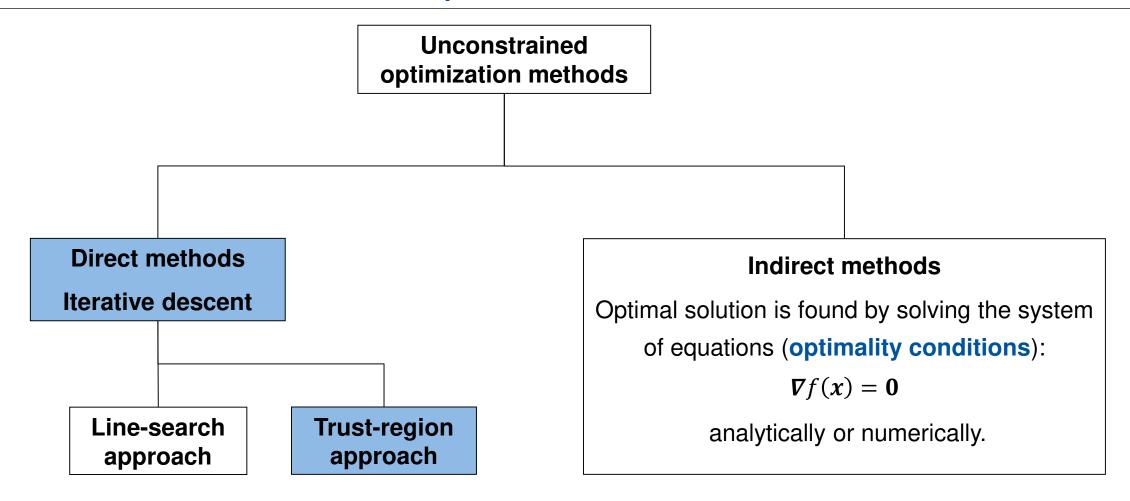
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Trust region method





# **Solution Methods for Unconstrained Optimization**





# **Trust Region Method (1)**

### Idea:

• Approximate f at  $x^{(k)}$  by the quadratic model function  $m^{(k)}$ :

$$m^{(k)}(\boldsymbol{p}) = f^{(k)} + \boldsymbol{g}^{(k)^T} \boldsymbol{p} + \frac{1}{2} \boldsymbol{p}^T \boldsymbol{B}^{(k)} \boldsymbol{p}$$
 where  $f^{(k)} = f(\boldsymbol{x}^{(k)})$ ,  $\boldsymbol{g}^{(k)} = \nabla f(\boldsymbol{x}^{(k)})$  and  $\boldsymbol{B}^{(k)}$  is symmetric

- Taylor-series: the approximation error is small for small p
- For each iteration k = 0,1,... choose a trust region radius  $\Delta^{(k)}$
- Solve the minimization problem:  $\min_{\boldsymbol{p}} m^{(k)}(\boldsymbol{p})$  s. t.  $\|\boldsymbol{p}\| \le \Delta^{(k)}$  and set  $\boldsymbol{p}^{(k)}$  to the solution found
- Set  $x^{(k+1)} = x^{(k)} + p^{(k)}$





# **Trust Region Method (2)**

## How to update the radius $\Delta^{(k)}$ ?

• Compare the agreement between the model function  $m_k$  and the objective function f at the previous iterations. Define contraction rate  $\rho_k$  as:

$$\rho_k = \frac{f(\mathbf{x}^{(k)}) - f(\mathbf{x}^{(k)} + \mathbf{p}^{(k)})}{m^{(k)}(\mathbf{0}) - m^{(k)}(\mathbf{p}^{(k)})} = \frac{\text{actual reduction}}{\text{predicted reduction}}$$

As  $m_k$  is minimized over a domain containing **0**:

$$m^{(k)}(\mathbf{0}) - m^{(k)}(\mathbf{p}^{(k)}) > 0$$

If  $\rho_k < 0$ ! reject this step (ascent)

If  $\rho_k \approx 1$ ! increase the radius: good agreement

If  $\rho_k \approx 0$ ! decrease the radius: poor agreement





# **Trust Region Method (3)**

## Basic Algorithm:

choose 
$$\Delta^{(max)} > 0, \Delta^{(0)} \in (0, \Delta^{(max)})$$
 and  $\eta \in [0, \frac{1}{4})$ 

for 
$$k = 0,1,...$$

calculate direction  $p^{(k)}$ , contraction rate  $\rho_k$ 

if 
$$\rho_k < \frac{1}{4}$$
,  $\Delta^{(k+1)} = \frac{||\boldsymbol{p}^{(k)}||}{4}$ 

else if 
$$\rho_k > \frac{3}{4}$$
 and  $||p^{(k)}|| = \Delta^{(k)}, \Delta^{(k+1)} = \min(2\Delta^{(k)}, \Delta^{(max)})$ 

else 
$$\Delta^{(k+1)} = \Delta^{(k)}$$

if 
$$\rho_k > \eta$$
,  $x^{(k+1)} = x^{(k)} + p^{(k)}$ 

else 
$$x^{(k+1)} = x^{(k)}$$





## **Trust Region Method (3)**

### Remarks:

- $\Delta^{(k)}$  is increased only if  $||p^{(k)}||$  reaches the boundary of the domain.
- Strategies for the efficient solution of the minimization problem for  $p^{(k)}$ :
  - The Cauchy point: minimum along the steepest descent direction  $(-\boldsymbol{g}^{(k)})$ , slow
  - The *Dogleg method*: applicable when  $B^{(k)}$  is positive definite, fast (superlinear)
  - Steihaug's approach for large sparse matrices



# **The Dogleg Method**

### Idea:

- For a large  $\Delta^{(k)}$ : Newton step,  $\boldsymbol{p}^{(k)} = \boldsymbol{p}^B = -\boldsymbol{B}^{(k)^{-1}}\boldsymbol{g}^{(k)}$ . Where  $\boldsymbol{p}^B$  is the unconstrained minimum of  $m_k$ ,  $\|\boldsymbol{p}^B\| \leq \Delta^{(k)}$ .
- For a small  $\Delta^{(k)}$ : search the solution along the direction  $-\boldsymbol{g}^{(k)}$
- For an intermediate  $\Delta^{(k)}$ : additionally calculate  $p^U = -\frac{\boldsymbol{g}^{(k)^T}\boldsymbol{g}^{(k)}}{\boldsymbol{g}^{(k)}}\boldsymbol{g}^{(k)}$

Where  $p^U$  is the unconstrained minimum of  $m^{(k)}$  in the steepest descent direction.



$$\boldsymbol{p}^{U} \text{ and } \boldsymbol{p}^{B} \colon \boldsymbol{p}^{(k)}(\tau) = \begin{cases} \tau \boldsymbol{p}^{U} & 0 \leq \tau \leq 1 \\ \boldsymbol{p}^{U} + (\tau - 1)(\boldsymbol{p}^{B} - \boldsymbol{p}^{U}) & 1 \leq \tau \leq 2 \end{cases}$$
 with  $\|\boldsymbol{p}^{(k)}(\tau^{*})\| = \Delta^{(k)}$ .



 $\Delta^{(k)}$ 

 $p^{(k)}(\tau^*)$ 

 $x^{(k)}$ 

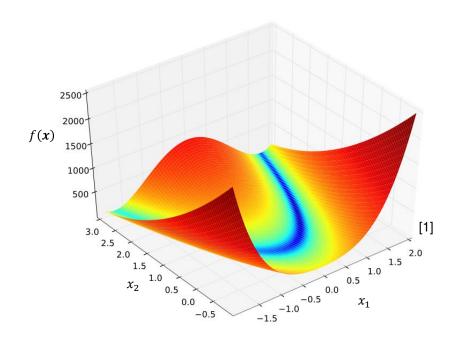


Dogleg-path

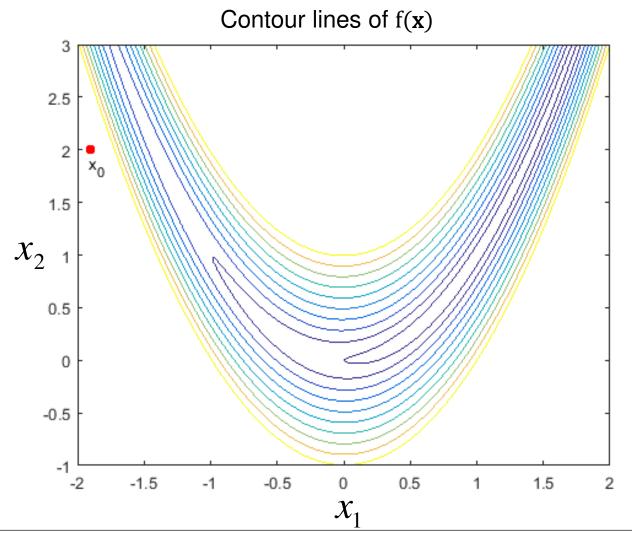
# **Illustration of Convergence – Rosenbrock Function**

$$\min_{\mathbf{x} \in R^2} f(\mathbf{x}) = 100(x_2 - x_1^2)^2 + (1 - x_1)^2$$

• Solution point is  $x = (1,1)^T$  - why?



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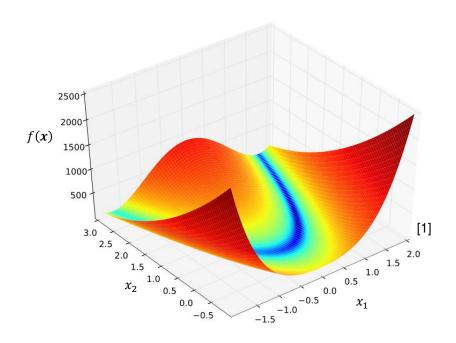




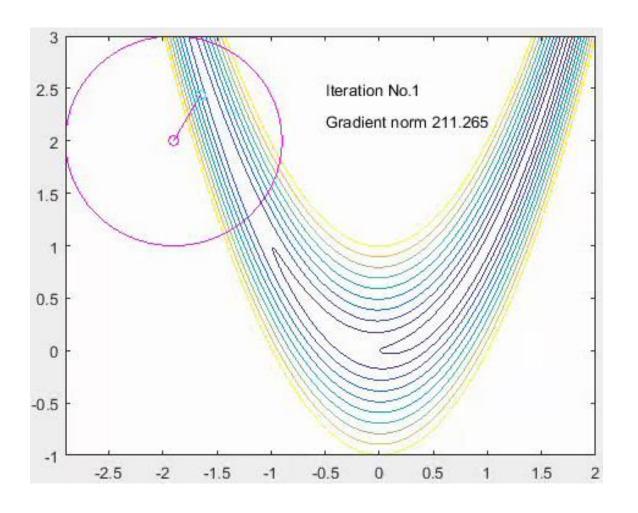
# Illustration of Convergence – BFGS (Quasi-Newton Method)

$$\min_{\mathbf{x} \in R^2} f(\mathbf{x}) = 100(x_2 - x_1^2)^2 + (1 - x_1)^2$$

Matlab trust-region (fminunc)



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## **Check Yourself**

- Explain trust region method.
- Which model problem is solved in trust-region methods?
- How is trust-region radius updated?



