Constant Optimizations

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- For example, searching for palindromic numbers can be done by iterating through integers in ascending order.
- Noticing that all palindromic integers are divisible by 11 allows for a constant optimization, where the time complexity remains the same.

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- (Python) using `sys.stdin` and `sys.stdout`, and only flushing when needed. Note that `print` and `input` may flush your output arbitrarily, possibly slowing your program unnecessarily.
- (Java): I do not know how to make it fast, if I/O is a bottleneck, use a different language maybe? Kattio.java exists but is not even comparable to C++ speed.

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- We will be examining a bit how data and instructions are processed by a CPU.
- Be aware that microbenchmarks such as the ones I made for the slides may not always mean the changes translate perfectly into efficiency for real world code.
- The intent is to teach you the ideas. Do not assume that doing things manually guarantees speed increases! Try it out!

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- Modern CPUs tend to separate these two and treat them differently.
- Recently referenced, or used, data and instructions tend to be reused.
- Temporal locality: A recently accessed (memory) address is likely to be accessed in the near future.
- Spatial locality: An address nearby a recently accessed address is likely to be accessed in the near future.

Spotting Locality

```
int sum(int arr[N]) {
    int result = 0;
    for (int i = 0; i < N; i++) {
        result += a[i];
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- Here the same instructions are used in each iteration of a short loop: temporal locality of instructions
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- Here instructions are processed sequentially: spatial locality of instructions

Summing by columns

```c
int sum_by_col(int arr[N][M]){
    int result = 0;
    for (int j = 0; j < M; j++) {
        for (int i = 0; i < N; i++) {
            result += a[i][j];
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- Swapping the order of our loops makes use of spatial locality.
- That way the CPU cache will be utilized better.

Summing by rows

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int sum_by_row(int arr[N][M]){
    int result = 0;
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A Scheduling Problem - From Errichto on CF

Problem description
There are N workers, where $1 \leq N \leq 5\,000$. There is a 30 day window for a two man group project.

Each worker is either available or unavailable each day. You are given the list of days on which each worker is available. The project can only be worked on if both group members are available. You may assume all workers are equally competent, so you only want to maximize the number of days they work together.

What is the best pair of workers to select?

A Scheduling Problem - Slow Solution

```cpp
vector<vector<int>> workers;
int intersection(int a, int b) {
    int i = 0, j = 0;
    int result = 0;
    while (i < N && j < N) {
        if (workers[a][i] == workers[b][j]) {
            result++, i++, j++;
        }
        else if (workers[a][i] < workers[b][j]) {
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int intersection(int a, int b) {
    int result = 0;
    int inter = workers[a] & workers[b];
    for (int i = 0; i < D; i++) {
        if (inter & 1) {
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- The loop still takes $D = 30$ steps, but we can do it faster.

A Scheduling Problem - Magic Solution

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    return __builtin_popcount(workers[a] & workers[b]);
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- Now our time complexity is $\mathcal{O}(N^2)$ and the code runs fast enough.
- But wait, there is more!

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- What if $D > 64$?

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- We must also keep in mind the word size w on our machines. Modern machines usually are 64-bit, so $w = 64$.
- We could make a size $\lceil \frac{D}{w} \rceil$ array using w-bit integers to store our set.
- Implementing this includes a lot of shifts and indexing and it is easy to make a mistake and get it wrong.

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- Bitsets use the method discussed previously, packing 8 booleans in each byte.

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- It is in essence a boolean array
- How much memory does `bool arr[1 « 30]` take?
- Despite booleans only requiring one bit, the data type `bool` is a byte.
- The array is therefore 1024^3 bytes or 1 Gibibyte.
- Bitsets use the method discussed previously, packing 8 booleans in each byte.
- A `bitset<1 « 30>` is therefore 1024^3 bits or 128 Mebibytes.

A Scheduling Problem - Now With Bitsets

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constexpr int D{ 365 };
constexpr int max_n{ 5'000 };
bitset<D> workers[max_n];
int intersection(int i, int j) {
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- Practice problem: Chef and Queries

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- Loops are the same, either you repeat or break from the loop.
- A higher number of iterations on average with less branching is often faster.

Ternary Sum

```
constexpr int N{ 100'000'000 };
int arr[N];
int regular_loop(int arr[N]) {
    int sm = 0;
    for (int i = 0; i < N; i++) {
        if (arr[i] % 3 == 0) sm -= arr[i];
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    }
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- This code adds up all the numbers in an array with a slight modification.

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- We can shave off a little by repeating ourselves.

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constexpr int N{ 100'000'000 };
int arr[N];
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    int sm = 0;
    for (int i = 0; i < N; i+=16) {
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        // ...
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Single Instruction, Multiple Data (SIMD)

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- See documentation for details on functions.

SIMD

```cpp
#include <x86intrin.h> // includes all SIMD intrinsics
constexpr int N{ 100'000'000 };
int arr[N];
int simd(int arr[N]) {
    __m128i sm = _mm_setzero_si128();
    __m128i one = _mm_set_epi32(1, 1, 1, 1);
    for (size_t i = 0; i < N; i+=4) {
        __m128i nums = _mm_loadu_si128((__m128i*) (arr+i));
        __m128i to_add = _mm_set_epi32(arr[i+3]%3, arr[i+2]%3, arr[i+1]%3, arr[i]%3);
        to_add = _mm_sub_epi32(to_add, one);
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Summary

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- Constant time improvements are sometimes easier than algorithmic improvements, and may prove sufficient.