

$$2. S = [B]S | \epsilon$$

$$B = (S)S | S$$

$$\text{First}(S) = \{ [, \epsilon \} \quad \text{Follow}(S) = \{), \$ \}$$

$$\text{First}(B) = \{ (, [, \epsilon \} \quad \text{Follow}(B) = \{],), \$ \}$$

$$1. E = AE' \quad E' = +E | \epsilon$$

$$A = BA' \quad A' = *A | \epsilon$$

$$E = '(E') | N$$

$$N = (0|1)^*$$

$$\text{First}(E) = \{ (, 0, 1 \} \quad \text{Follow}(E) = \{), \$ \}$$

$$\text{First}(E') = \{ +, \epsilon \} \quad \text{Follow}(E') = \{), \$ \}$$

$$\text{First}(A) = \{ (, 0, 1 \} \quad \text{Follow}(A) = \{ +,), \$ \}$$

$$\text{First}(A') = \{ *, \epsilon \} \quad \text{Follow}(A') = \{ +,), \$ \}$$

$$\text{First}(B) = \{ (, 0, 1 \} \quad \text{Follow}(B) = \{ *, +,), \$ \}$$

$$\text{First}(N) = \{ 0, 1 \} \quad \text{Follow}(N) = \{), +, *, \$ \}$$

В обеих грамматиках нет конфликтов, и
они не леворекурсивны \Rightarrow LL(1) грамматиками